



<https://hao-ai-lab.github.io/dsc204a-f25/>

# DSC 204A: Scalable Data Systems

## Fall 2025

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# Large Message

Communication Model:  $\alpha + n\beta, \beta = \frac{1}{B}$

- The second term dominates – we want to minimize the second term
  - We want to utilize the bandwidth as much as possible

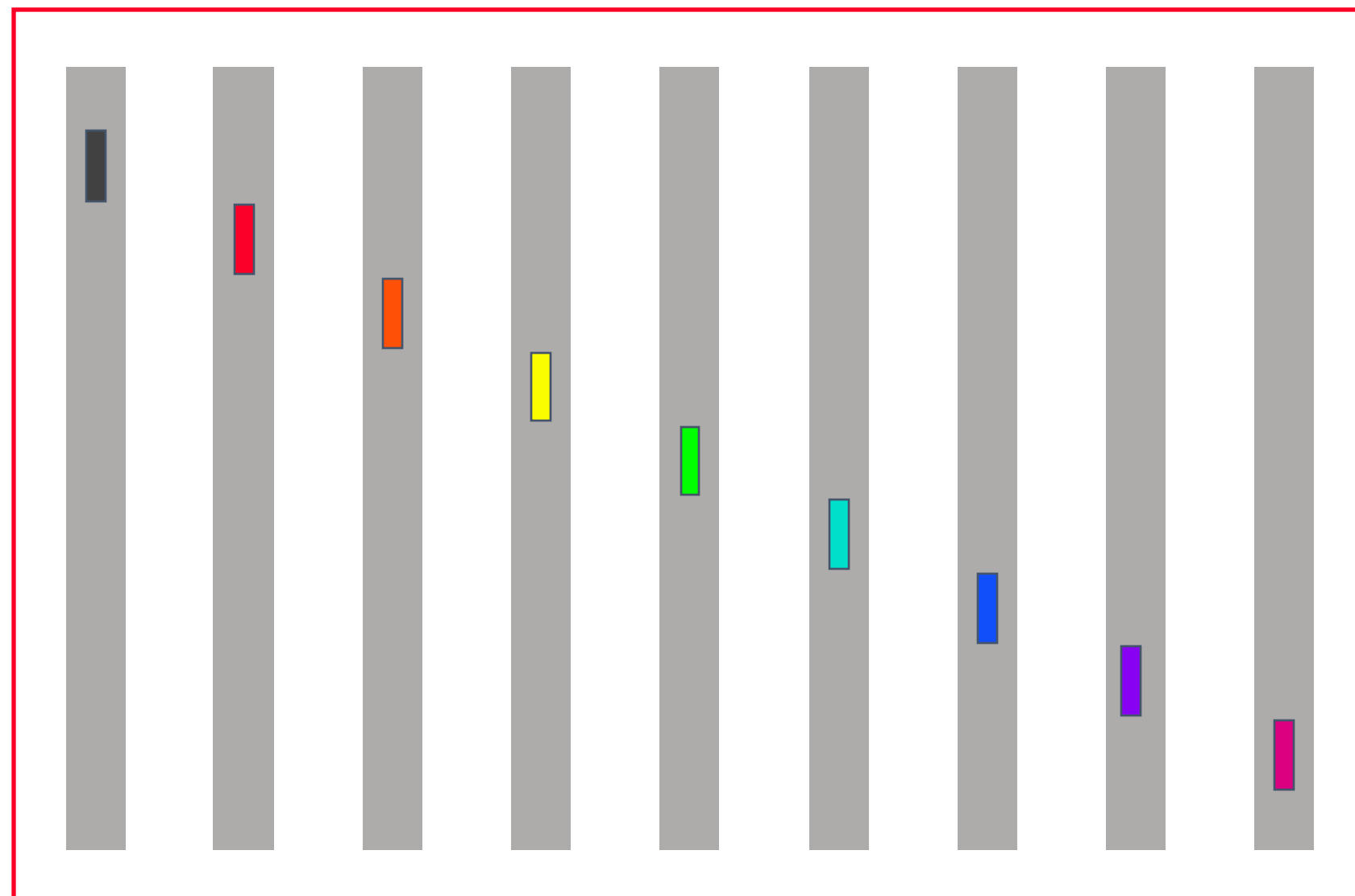
# General principles

Ring algorithm has the following advantages

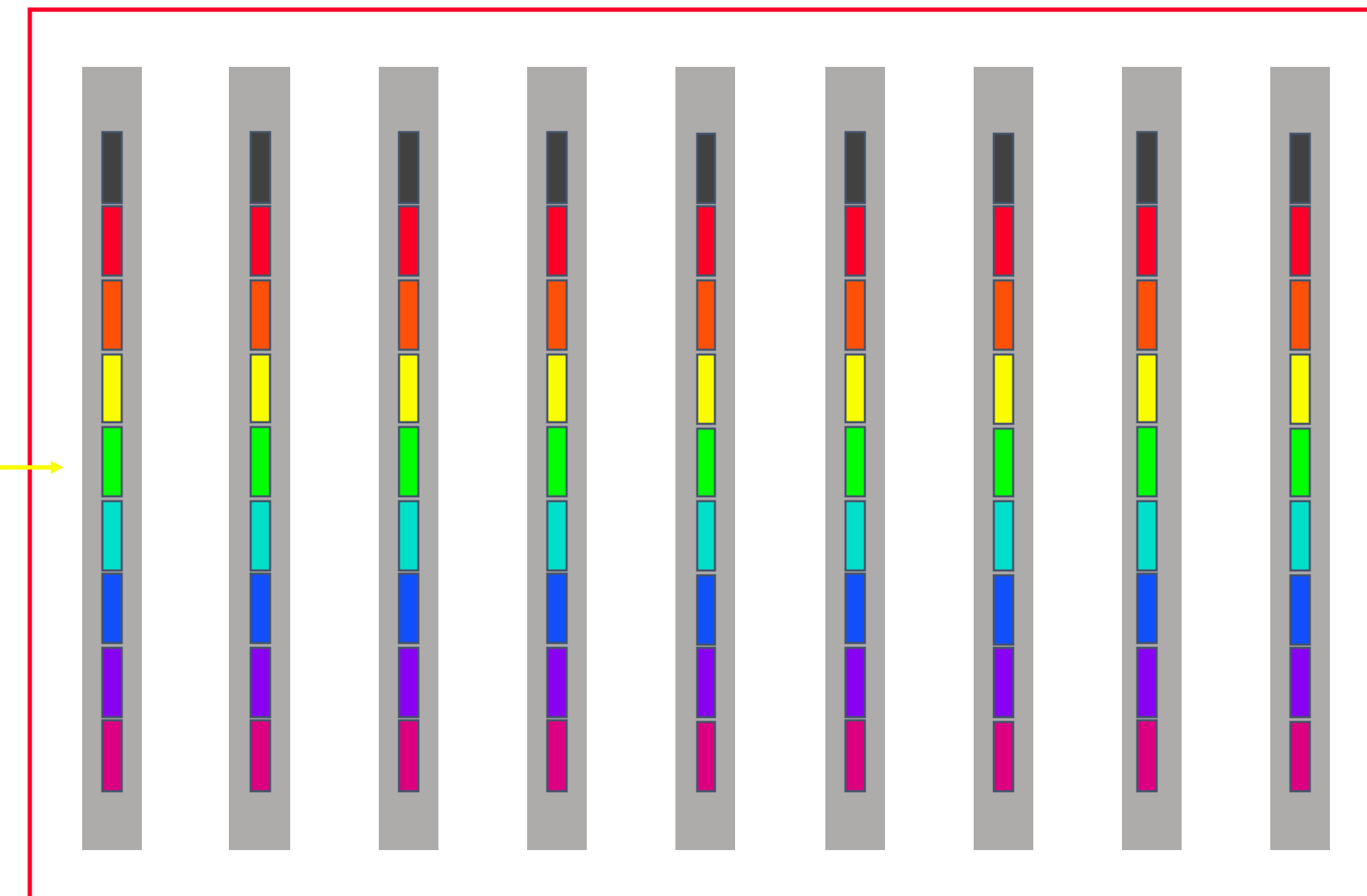
- Fully utilize the bandwidth (bandwidth optimal)
- implementation for arbitrary numbers of node

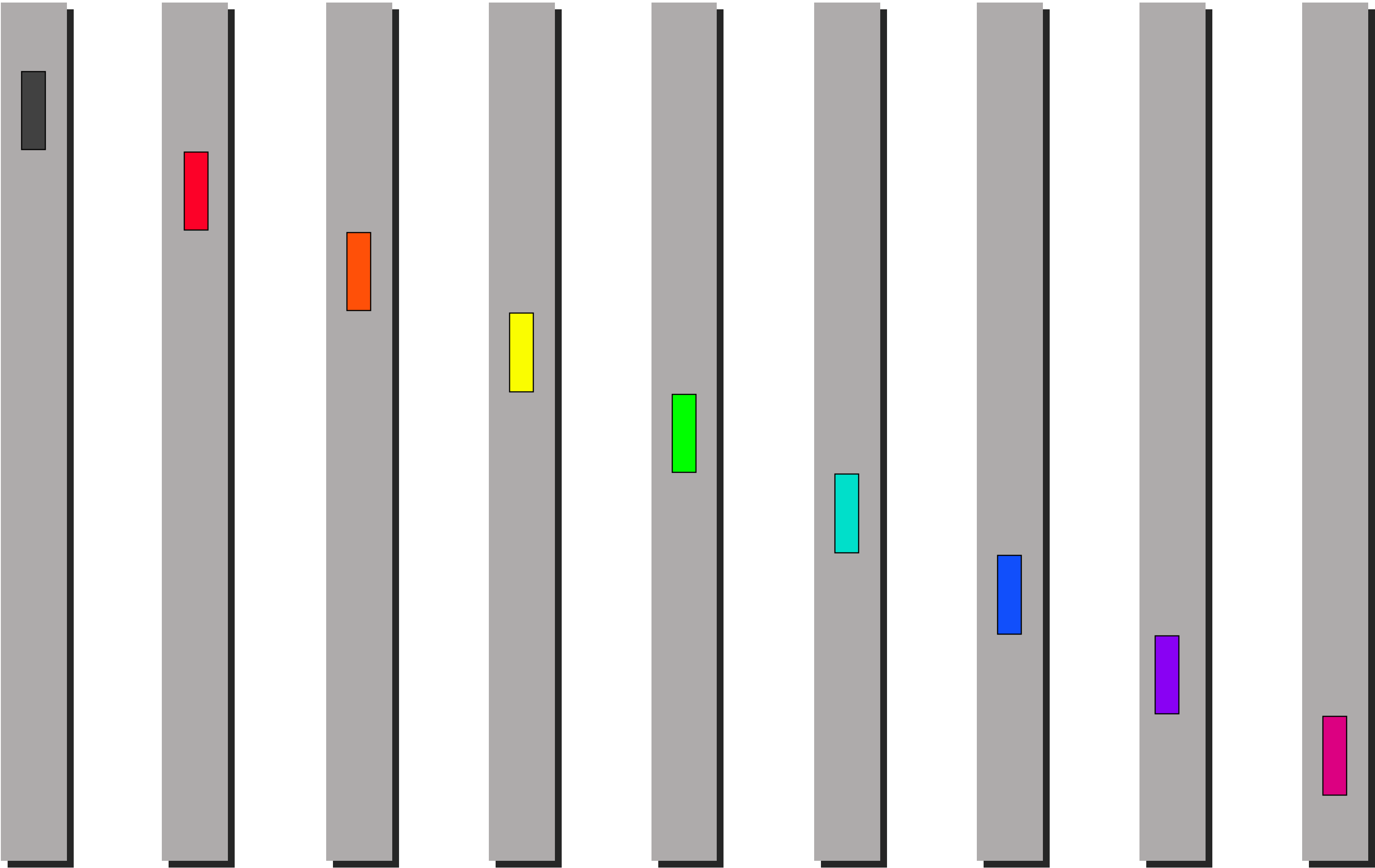
# Allgather

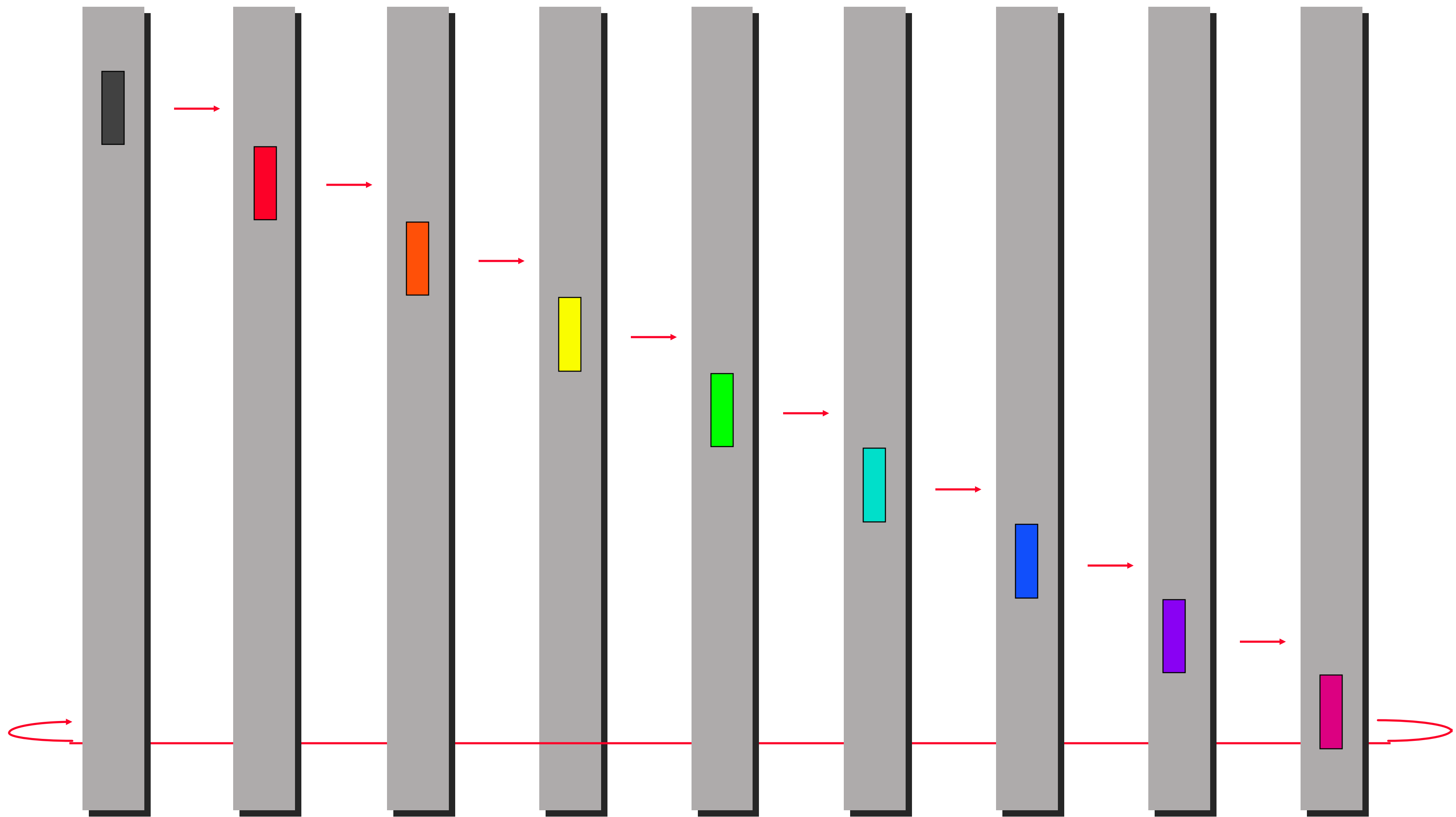
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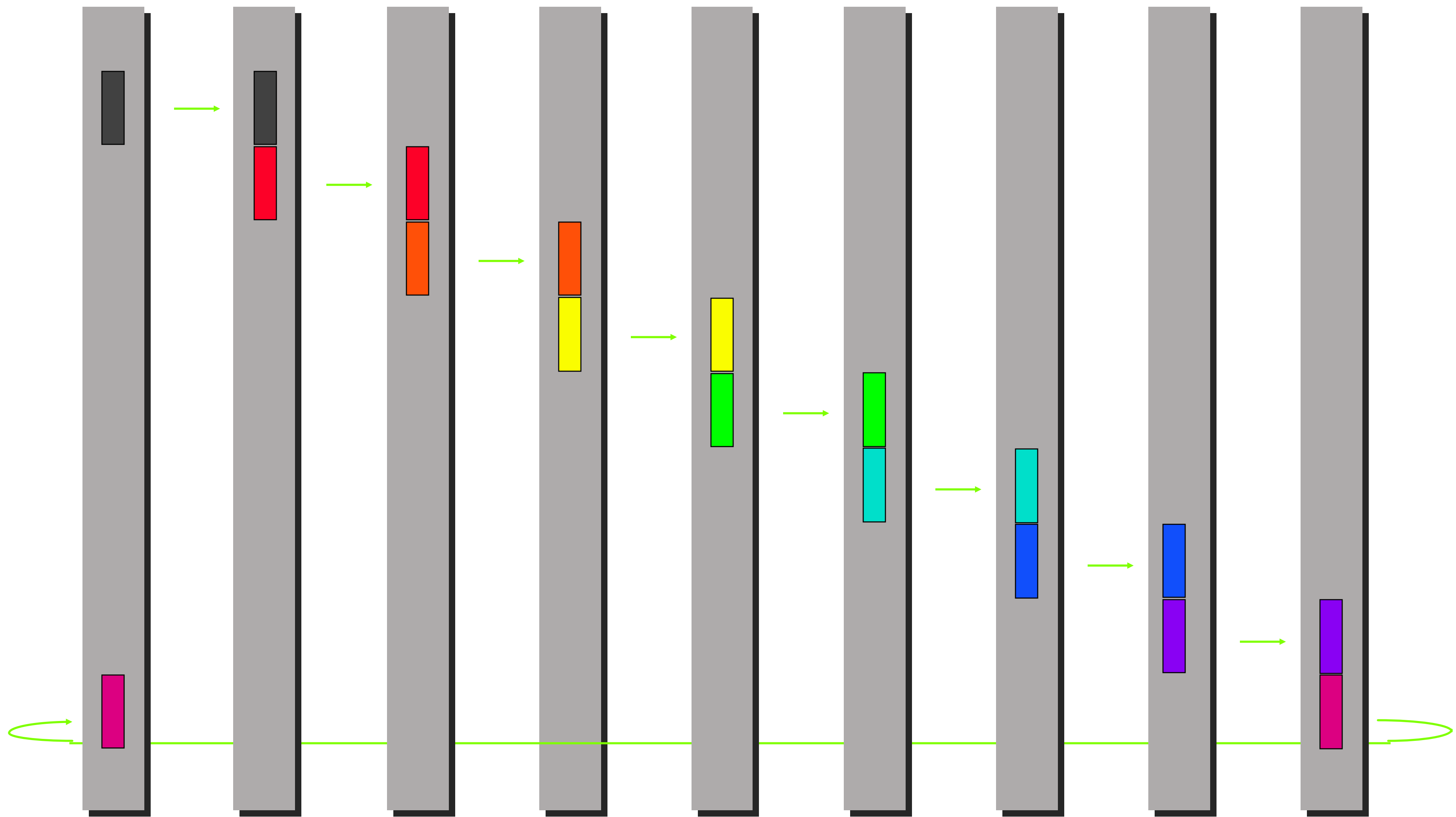


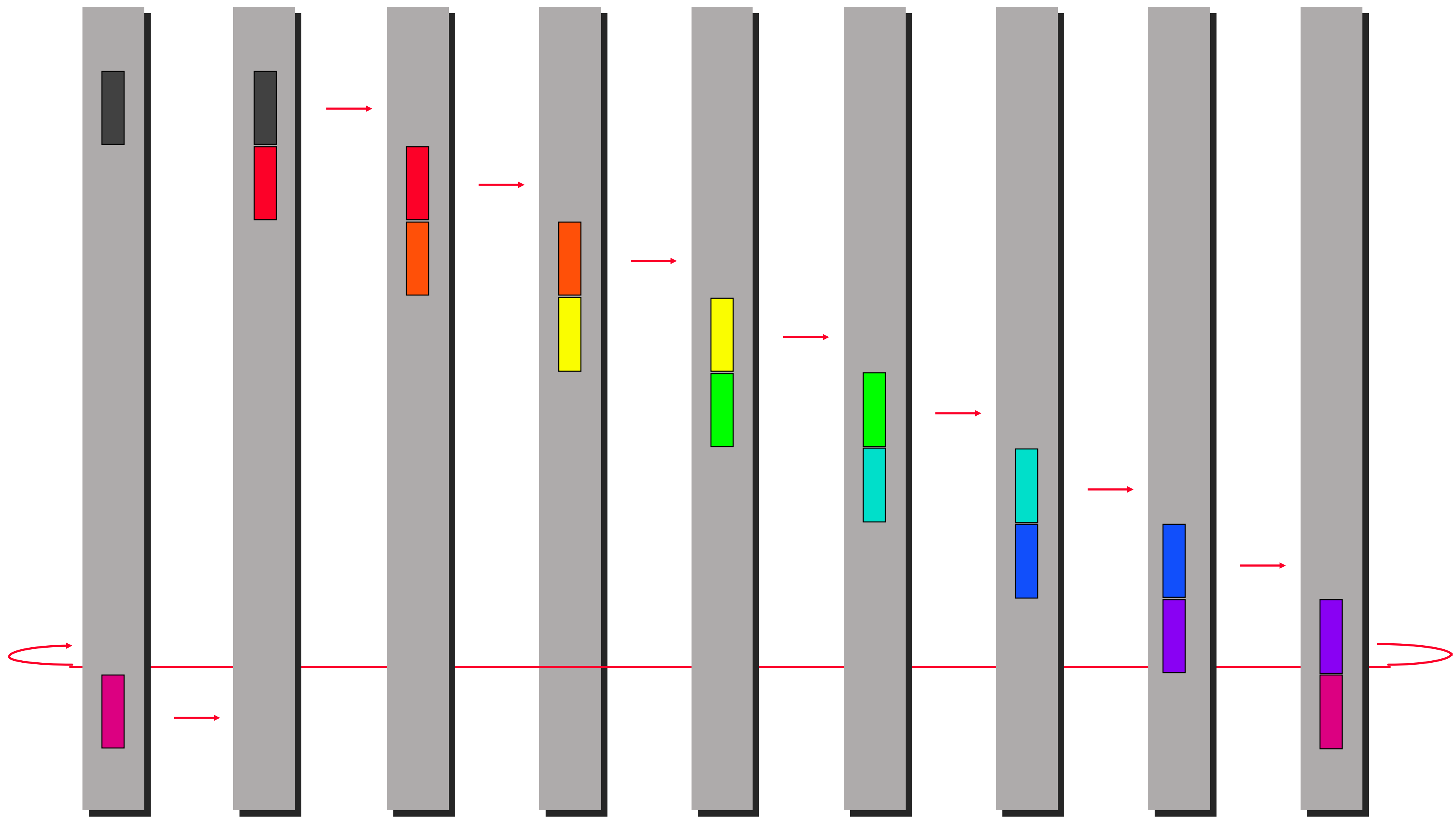
After



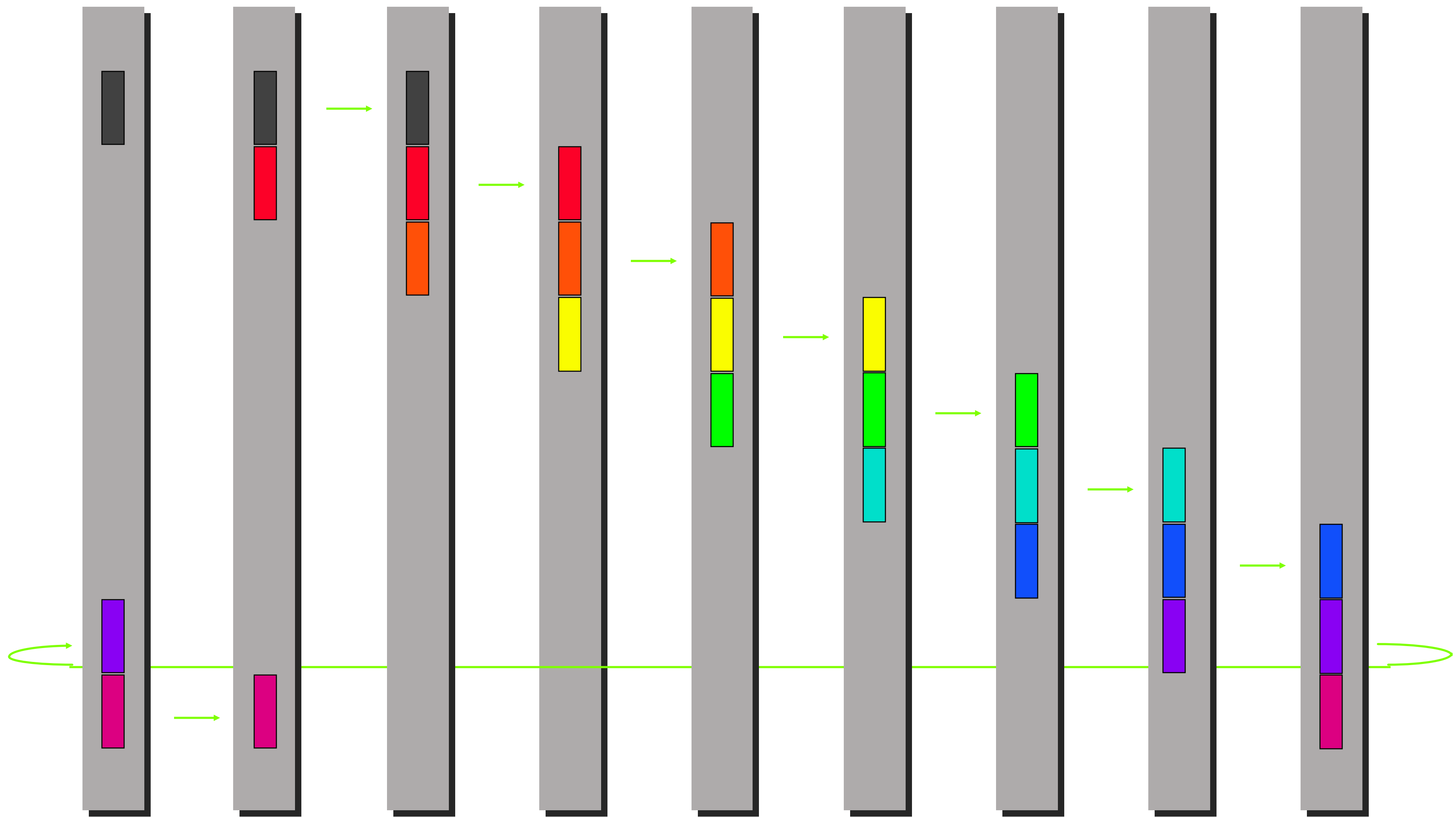


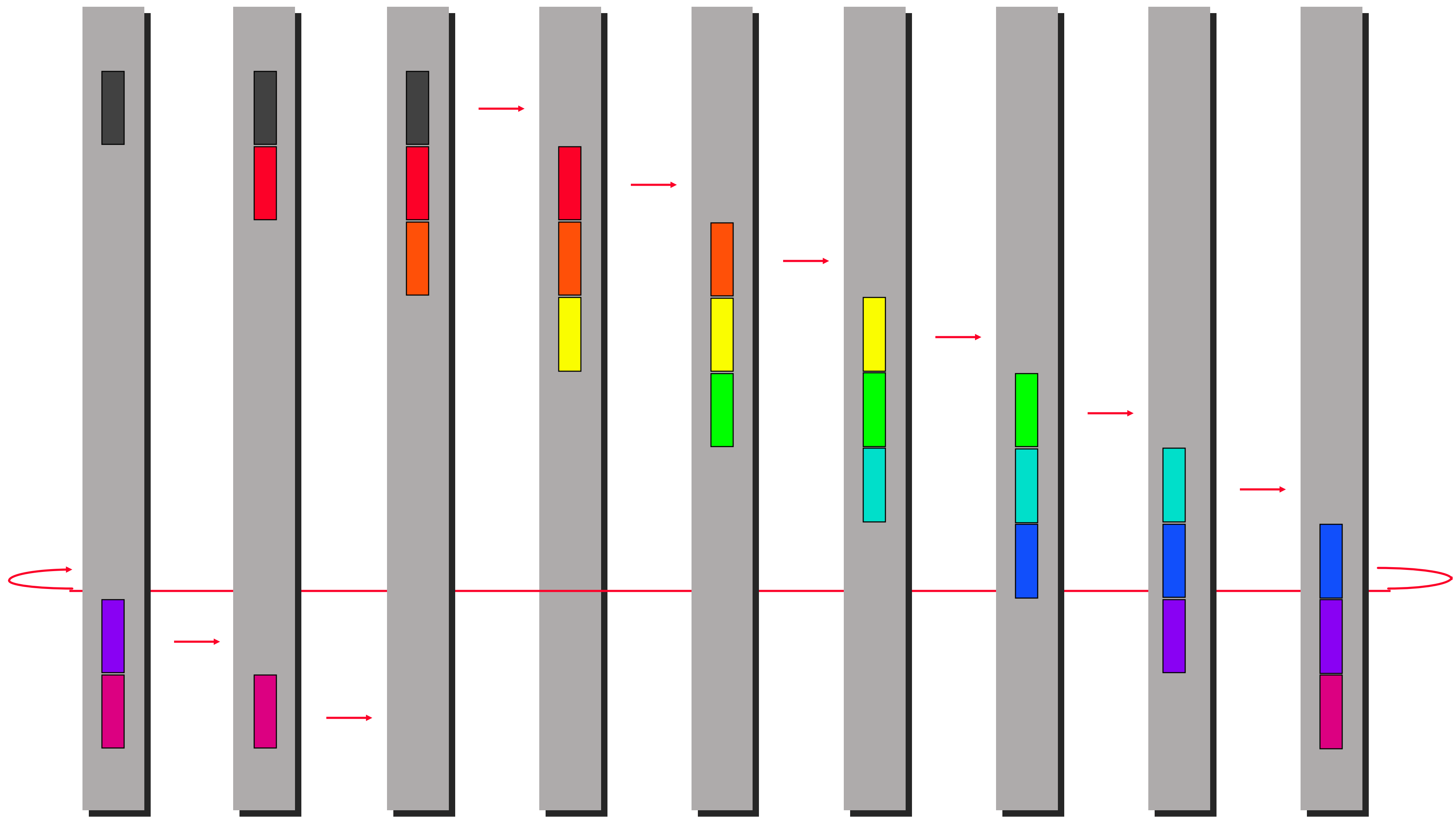


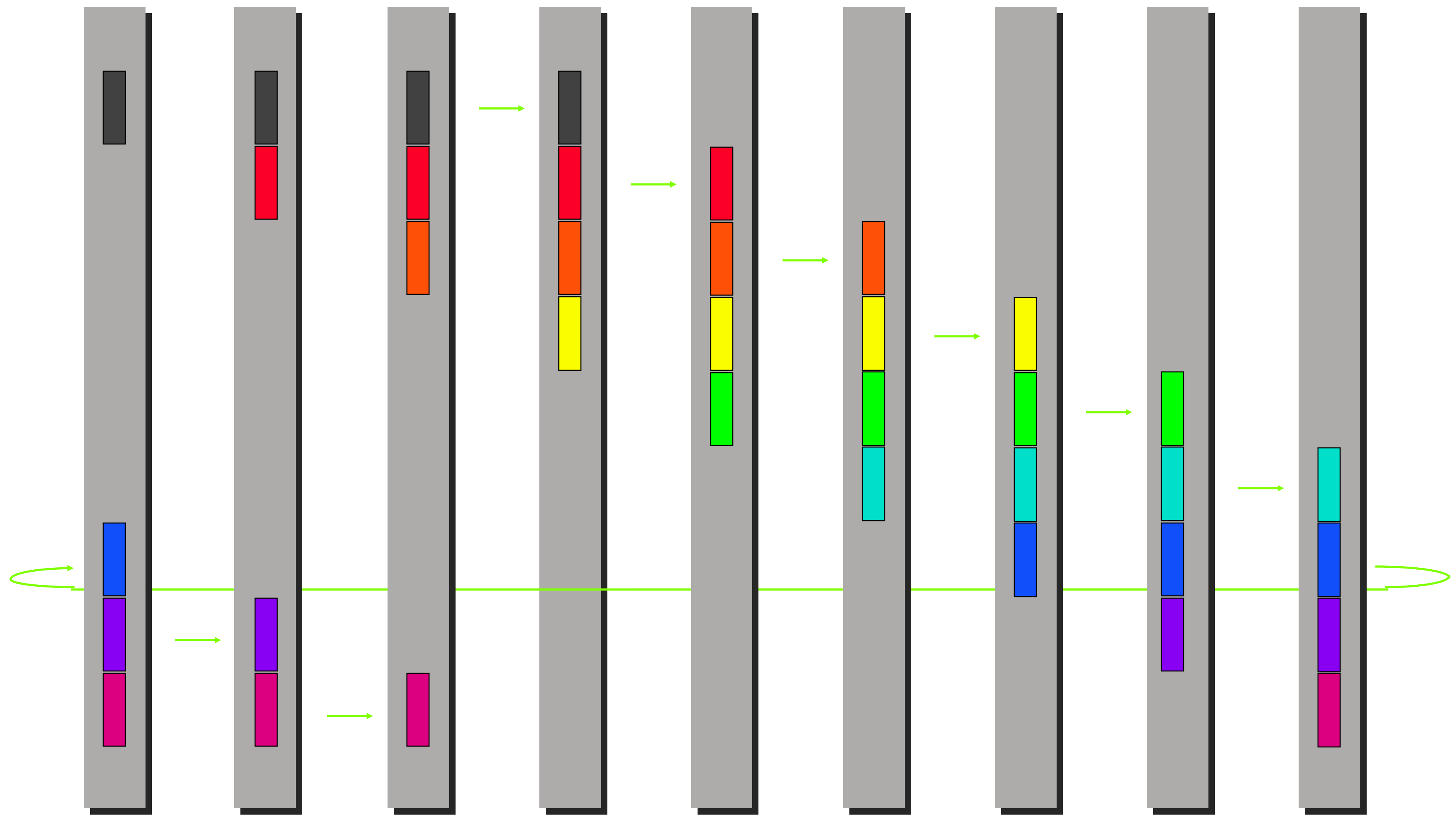


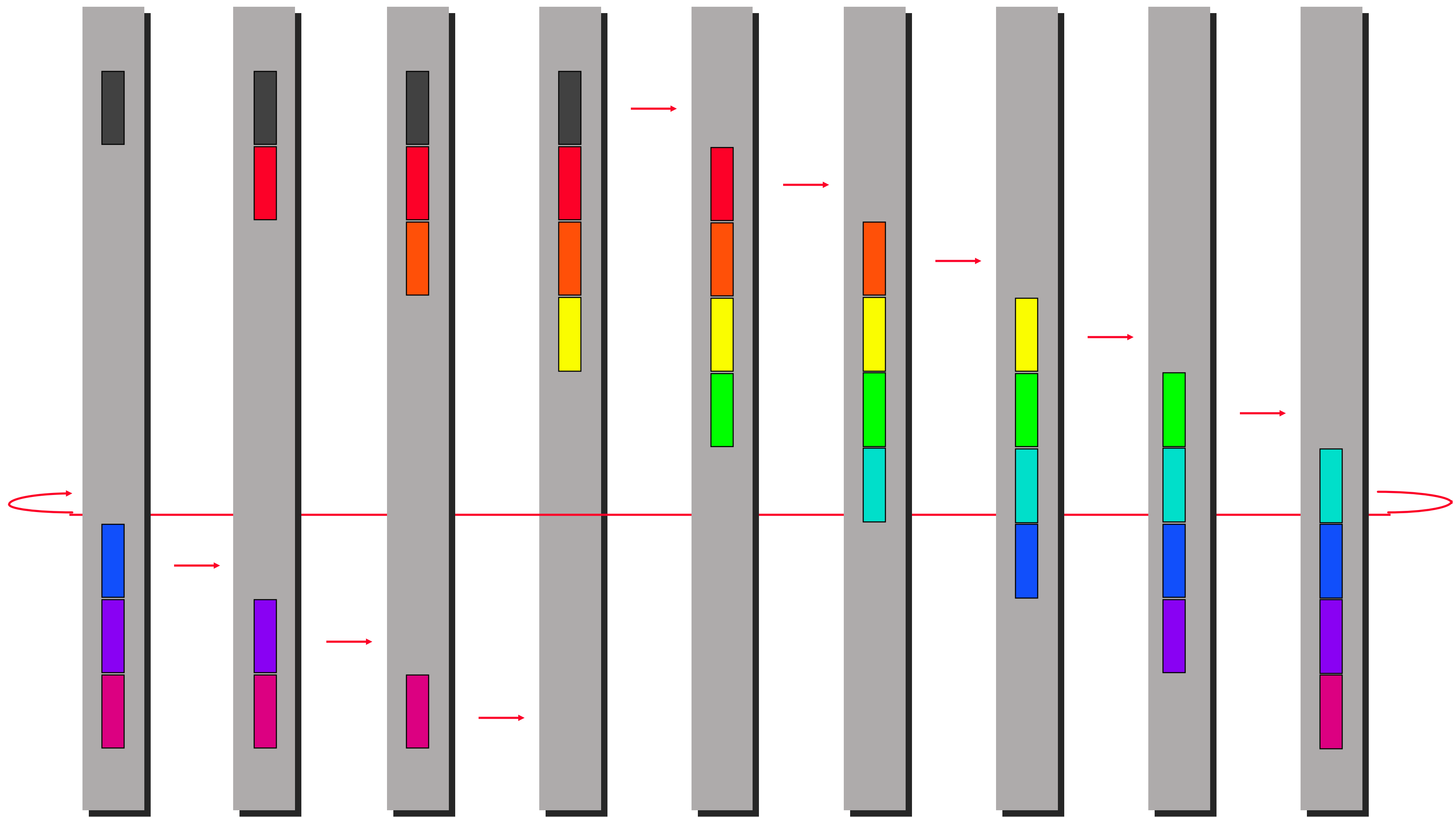


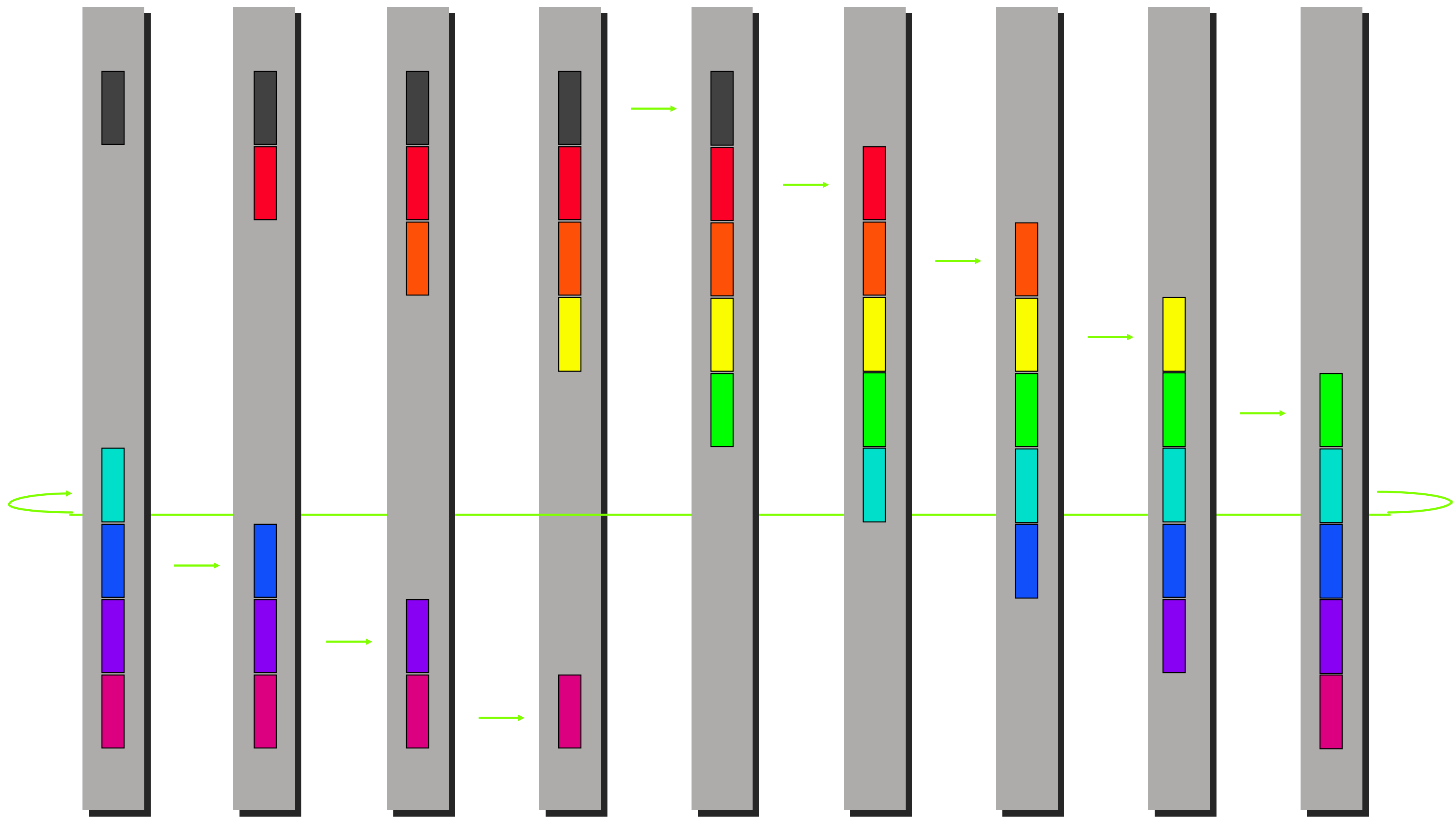


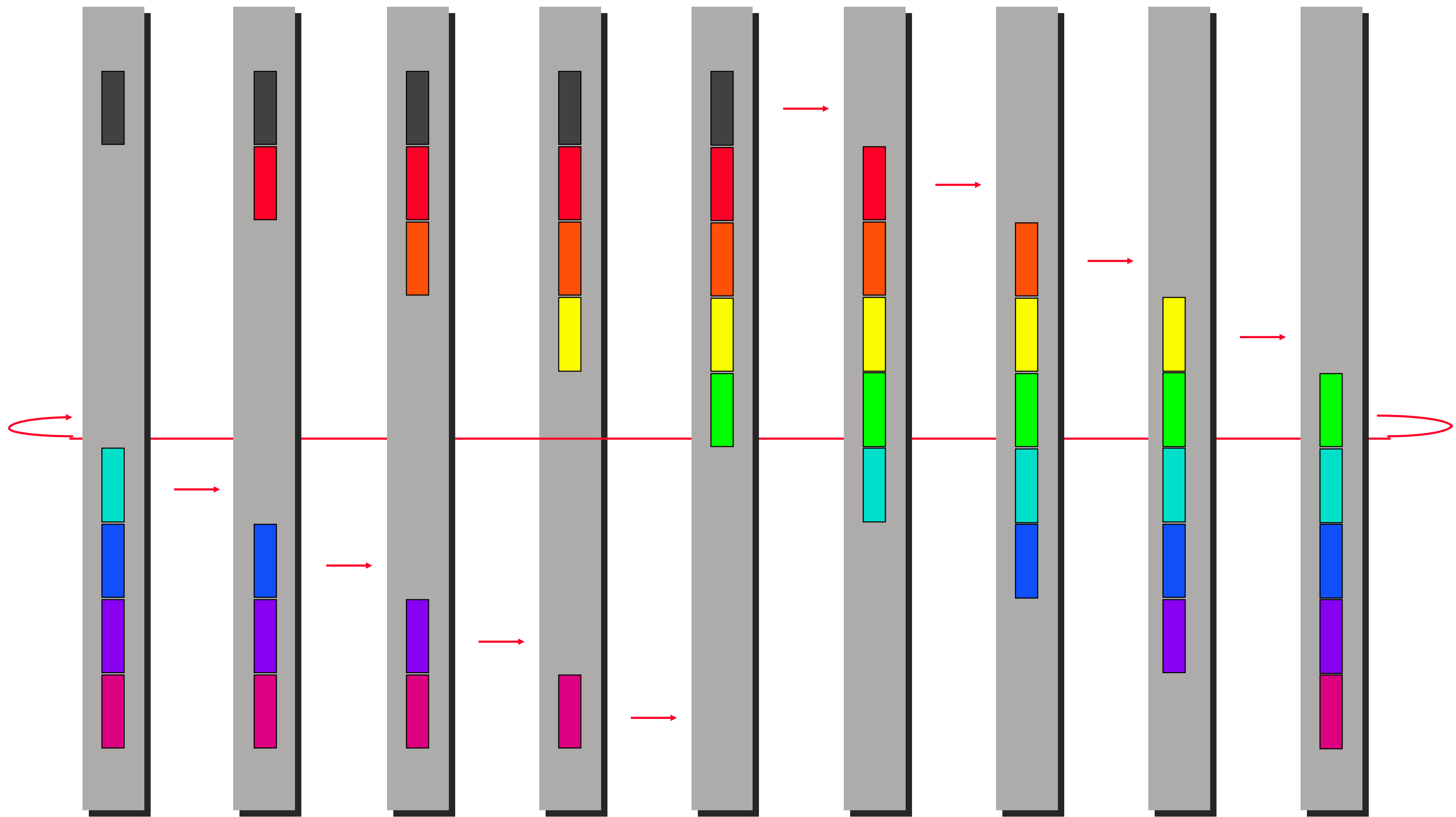


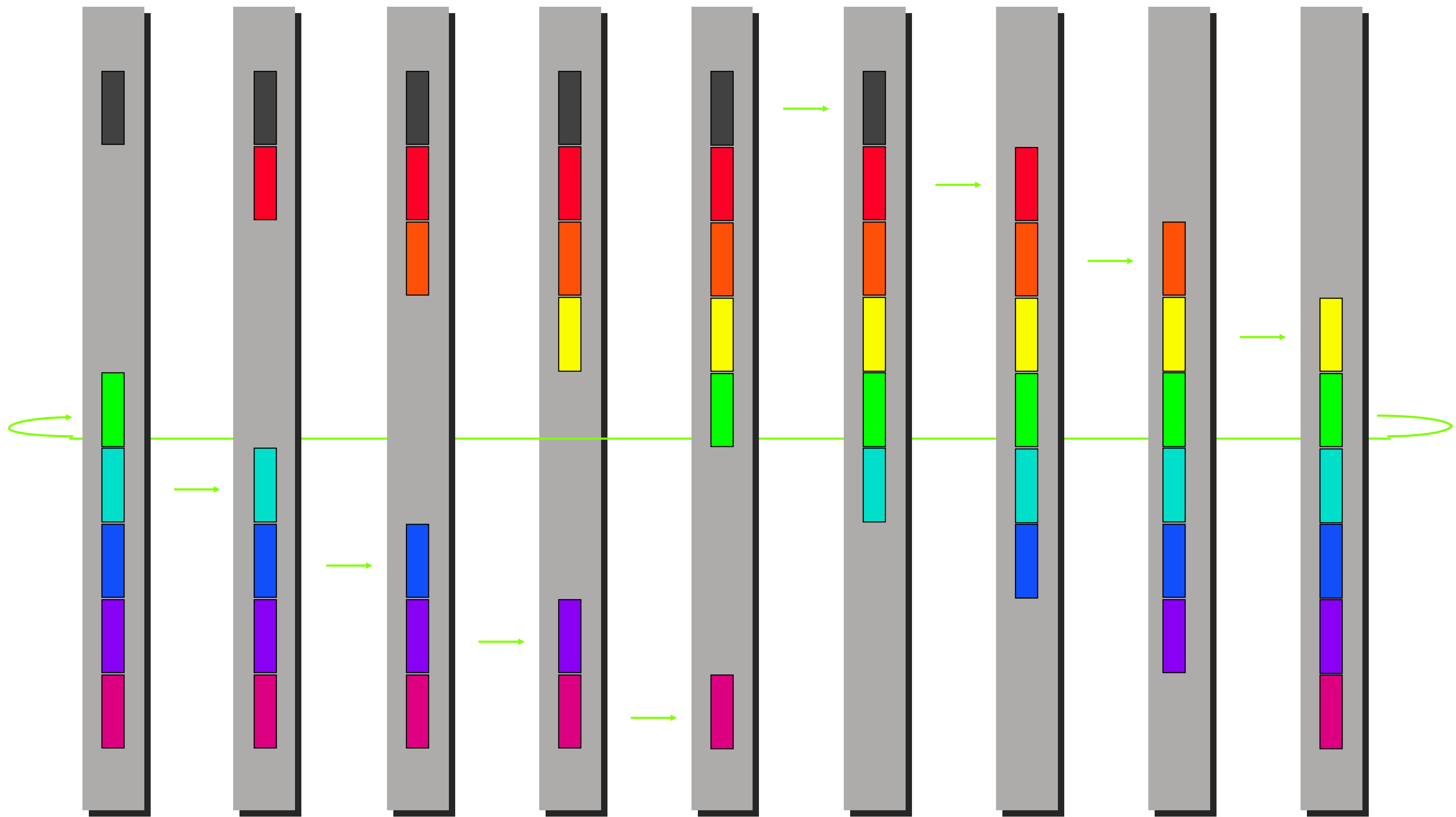


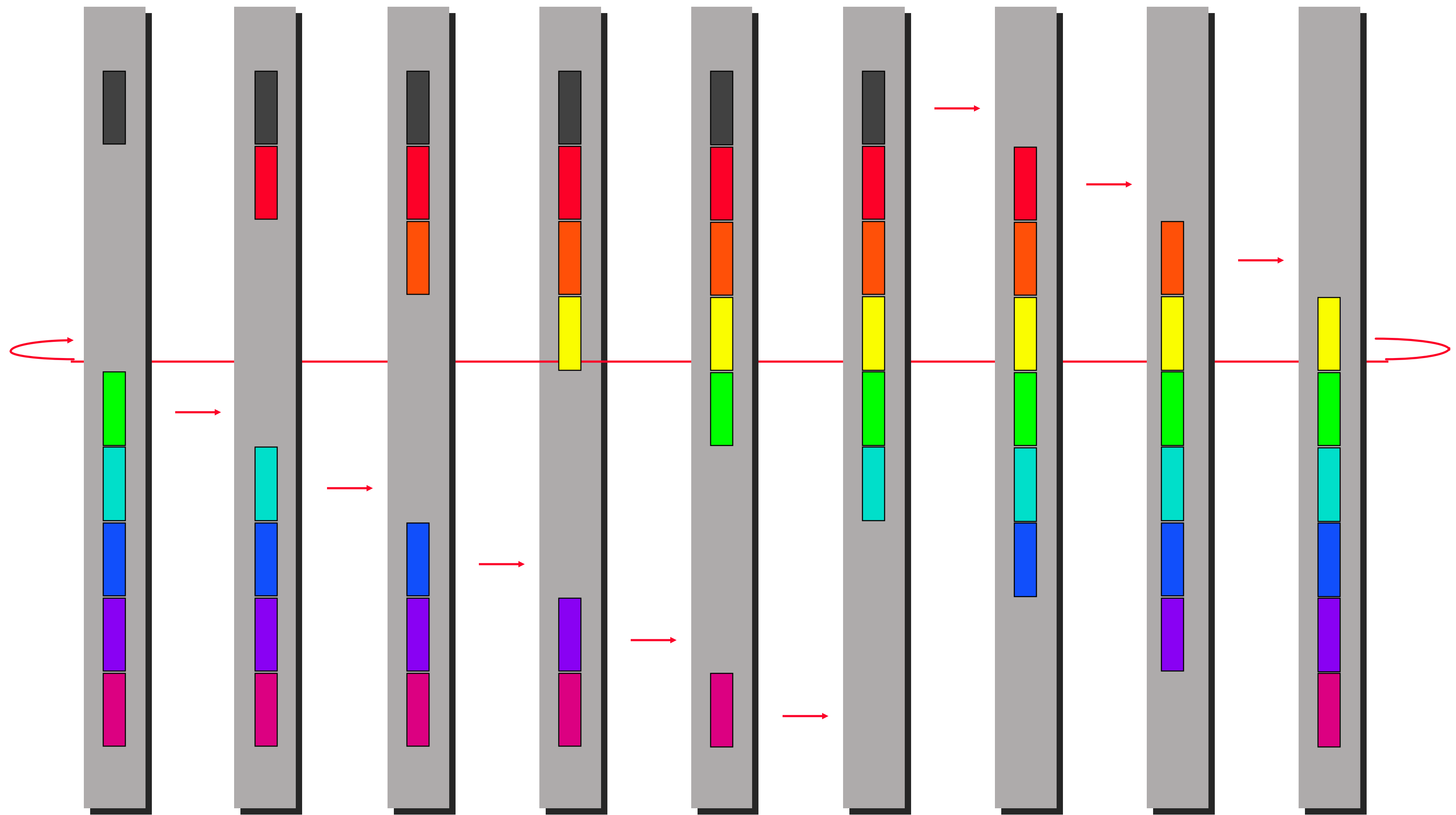




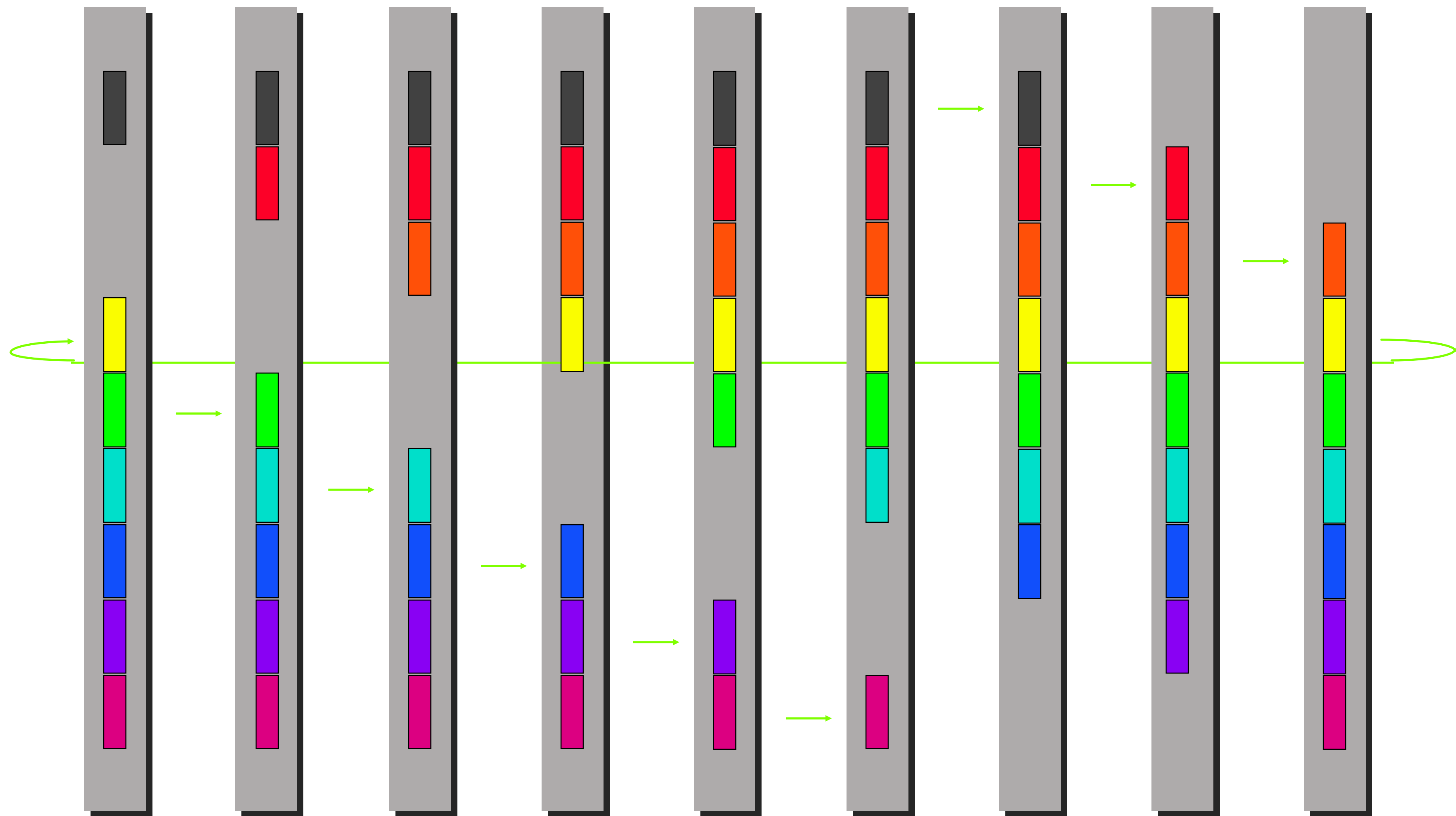


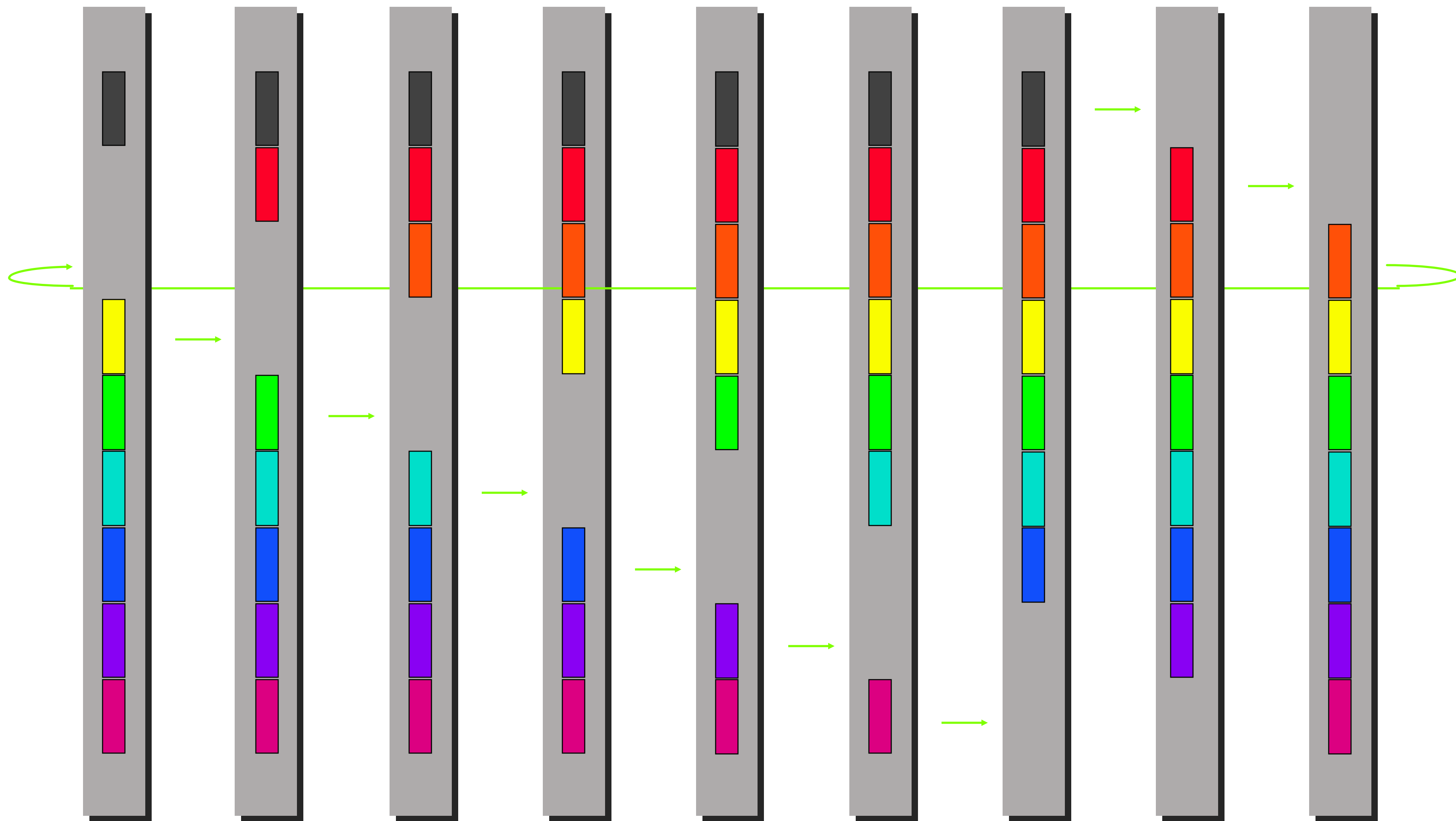


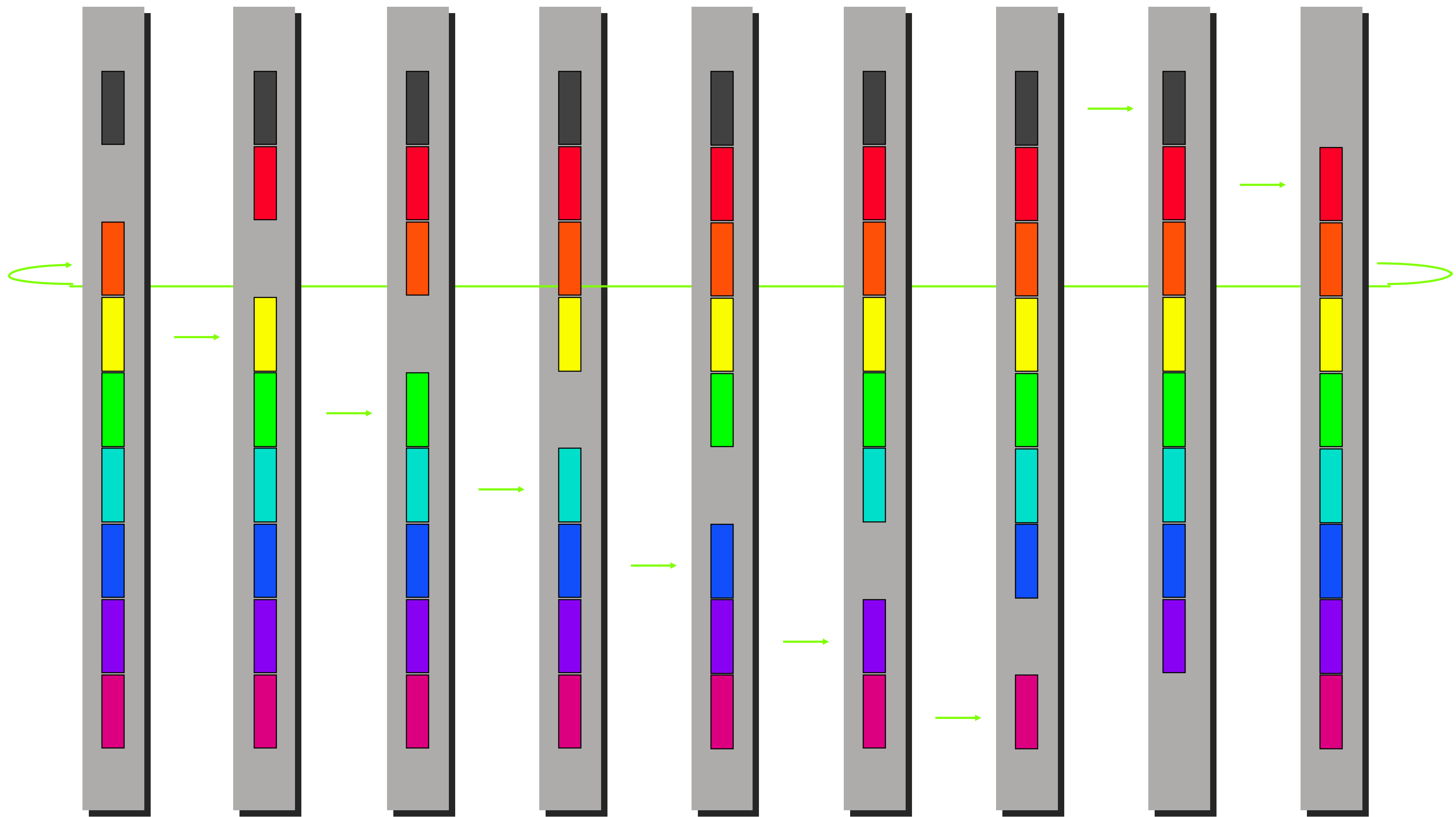


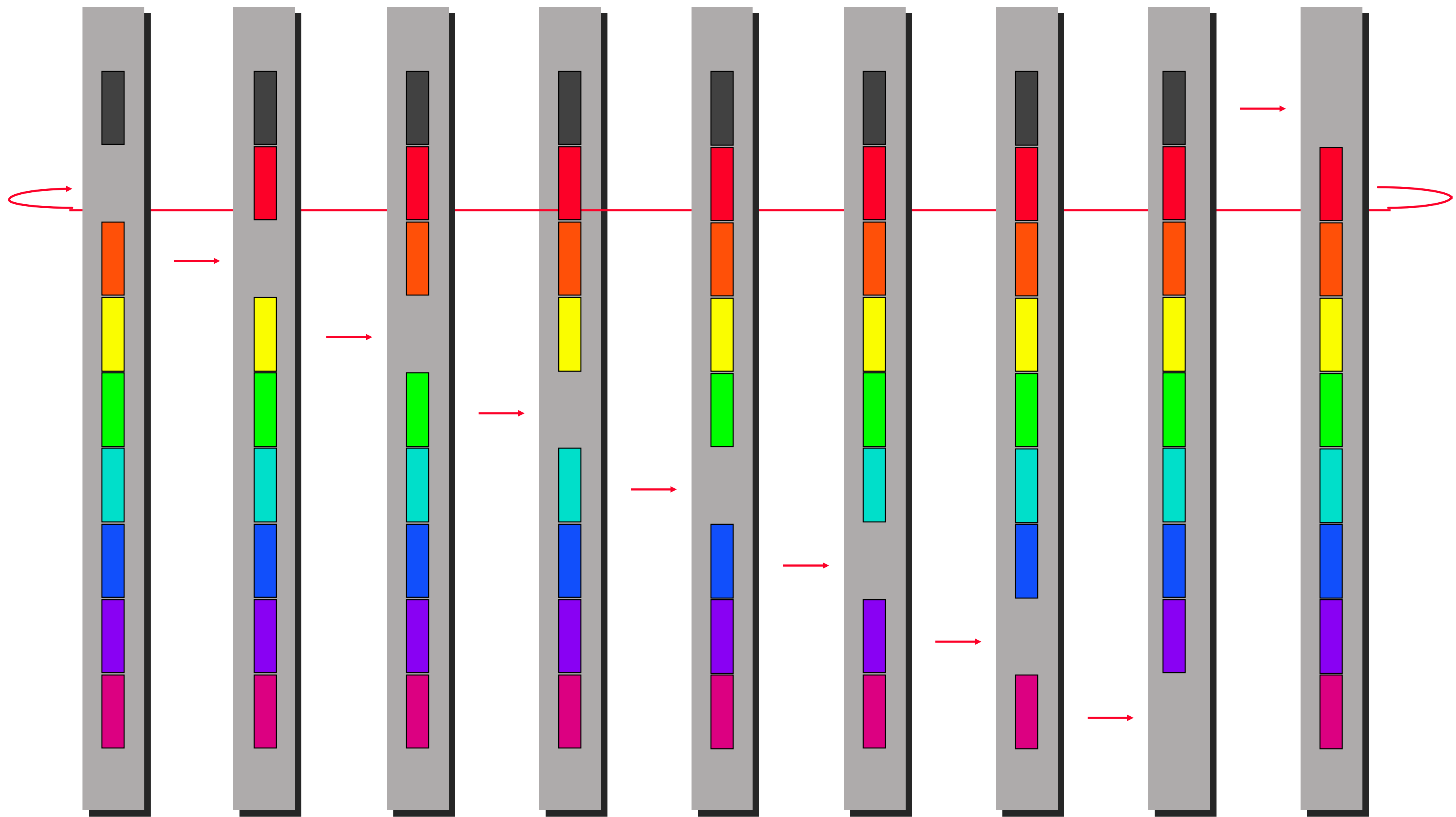


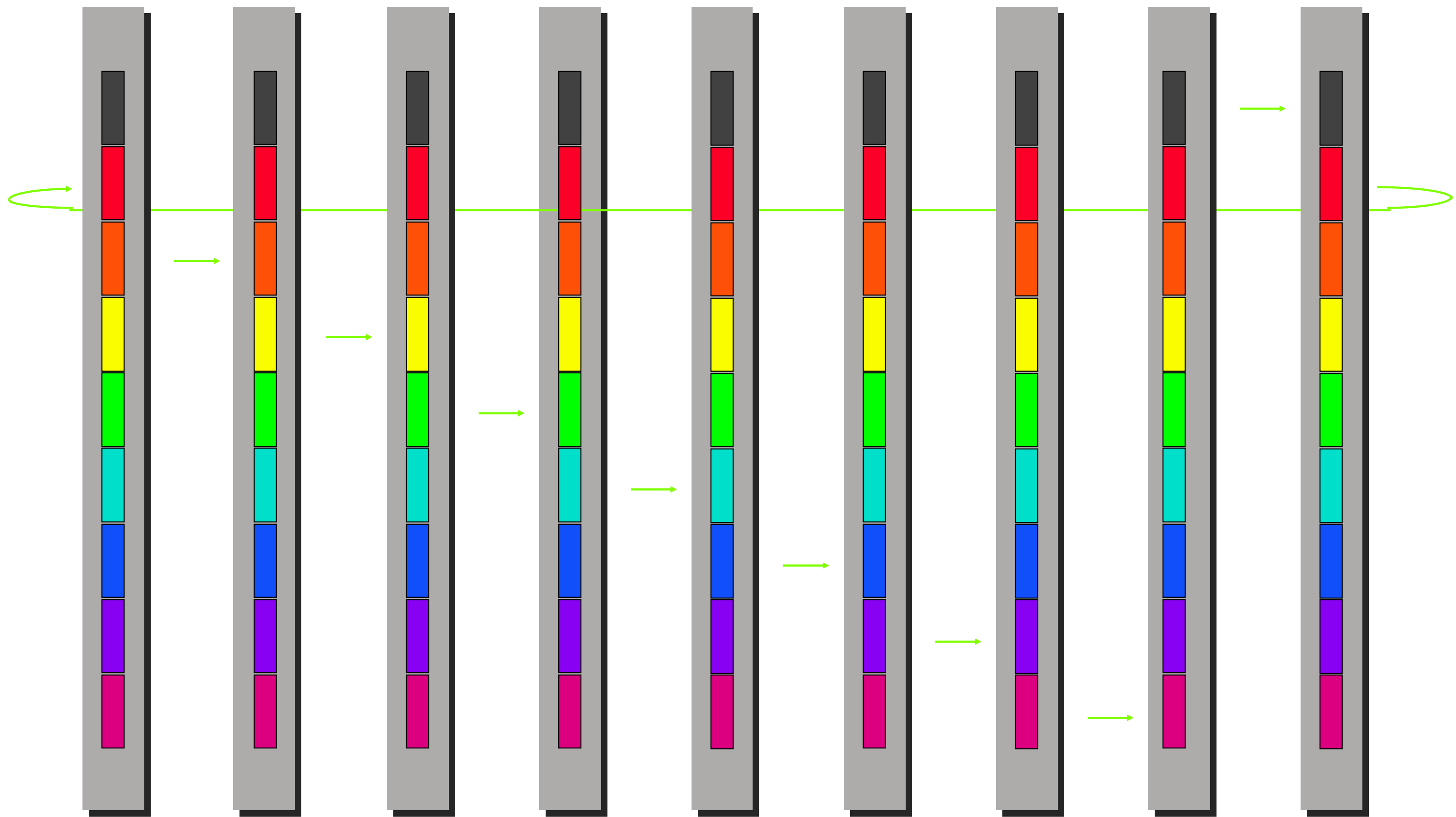


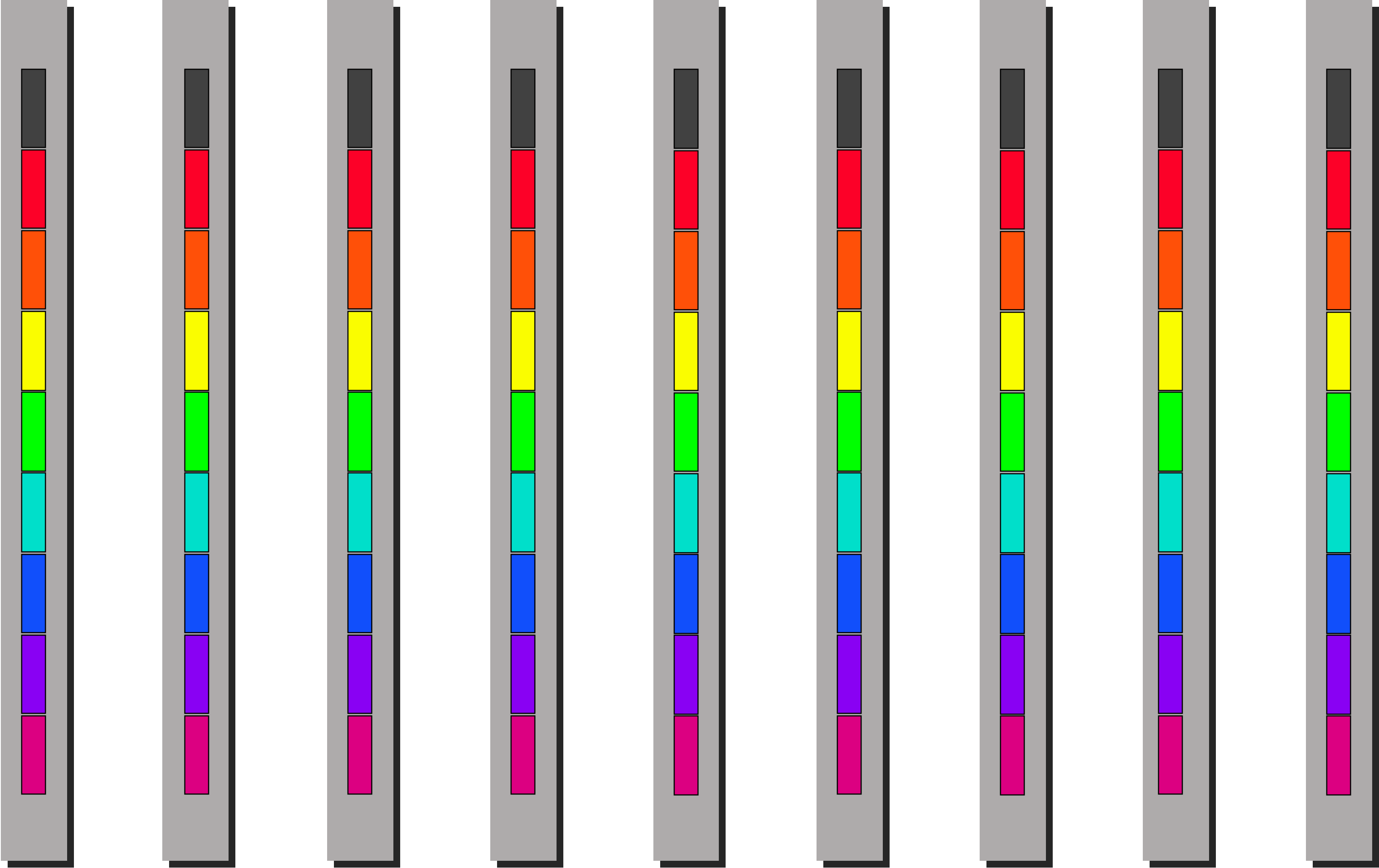




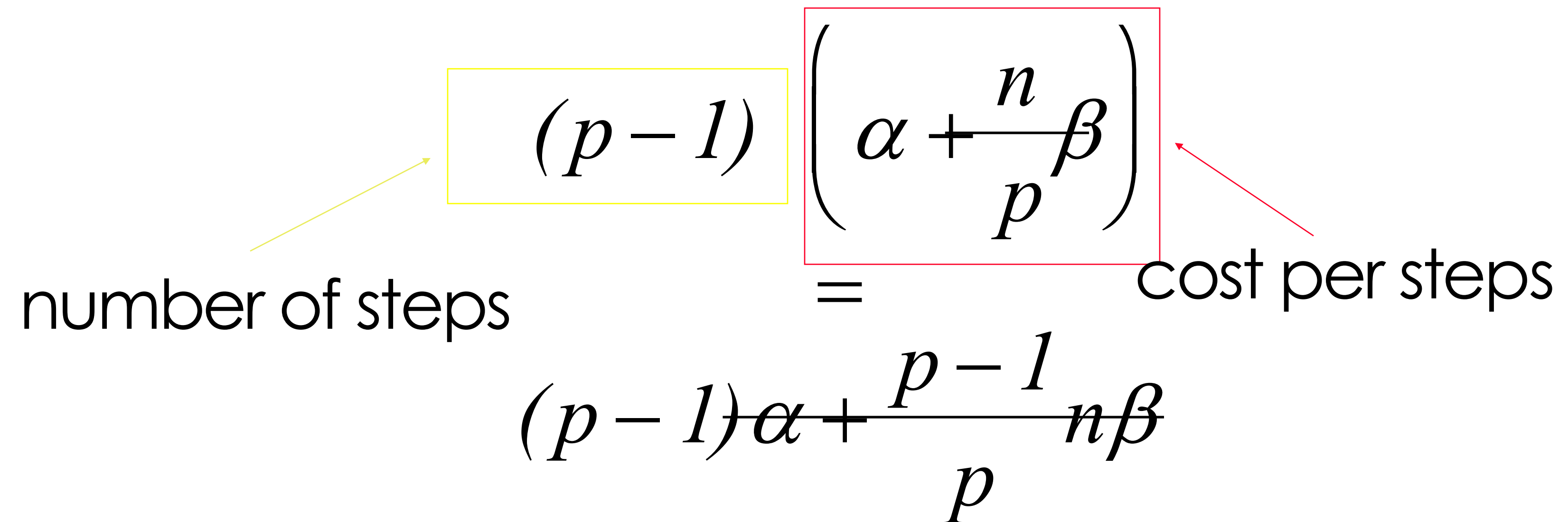








# Cost of bucket Allgather



The diagram illustrates the cost of bucket Allgather. It features a yellow box around the term  $(p-1)$  and a red box around the term  $\left(\alpha + \frac{n}{p}\beta\right)$ . A yellow arrow points from the text "number of steps" to the yellow box, and a red arrow points from the text "cost per steps" to the red box. Below these boxes, an equals sign is followed by the expanded formula  $(p-1)\alpha + \frac{p-1}{p}n\beta$ .

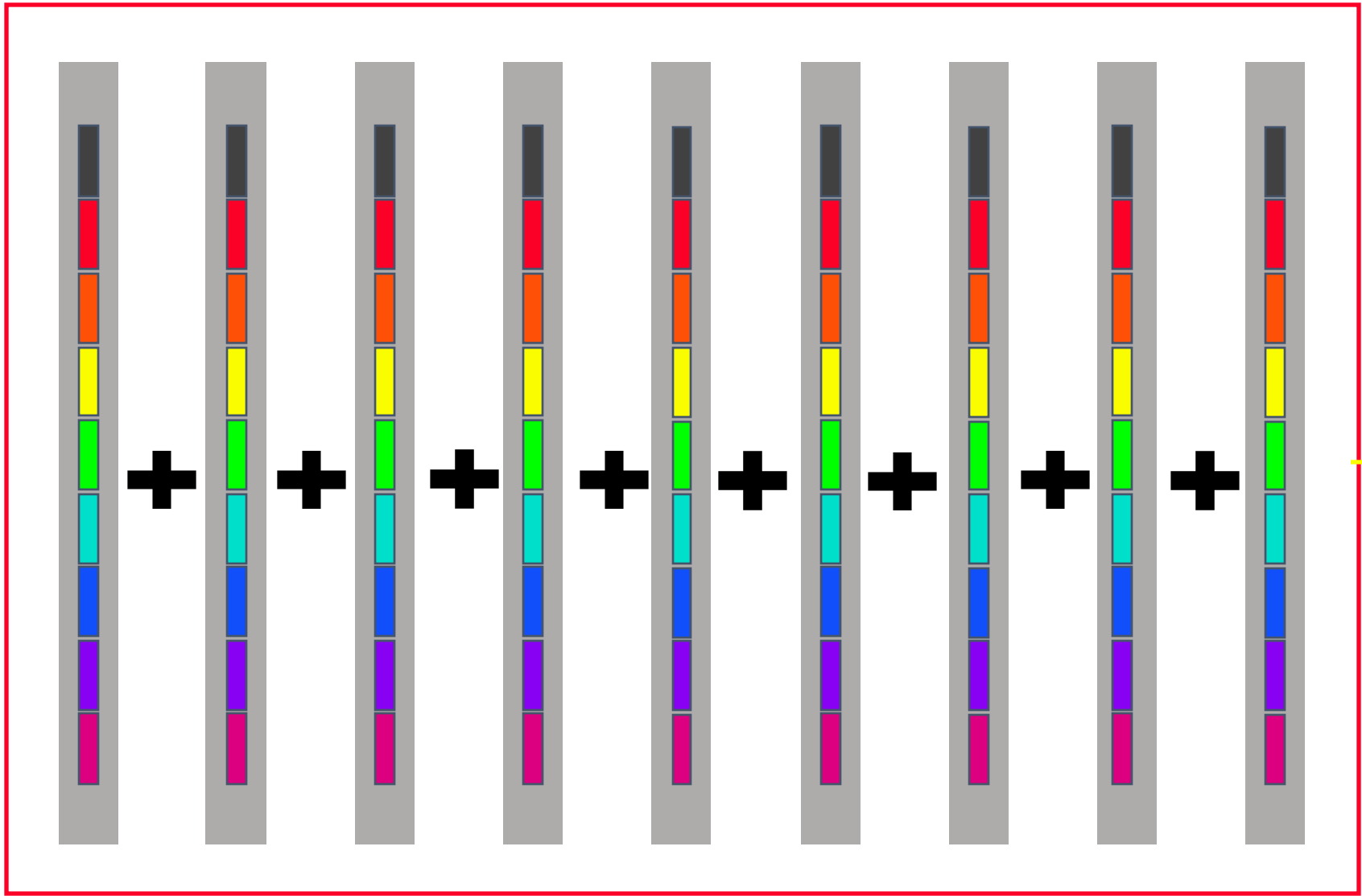
$$(p-1) \left( \alpha + \frac{n}{p}\beta \right) = (p-1)\alpha + \frac{p-1}{p}n\beta$$

number of steps

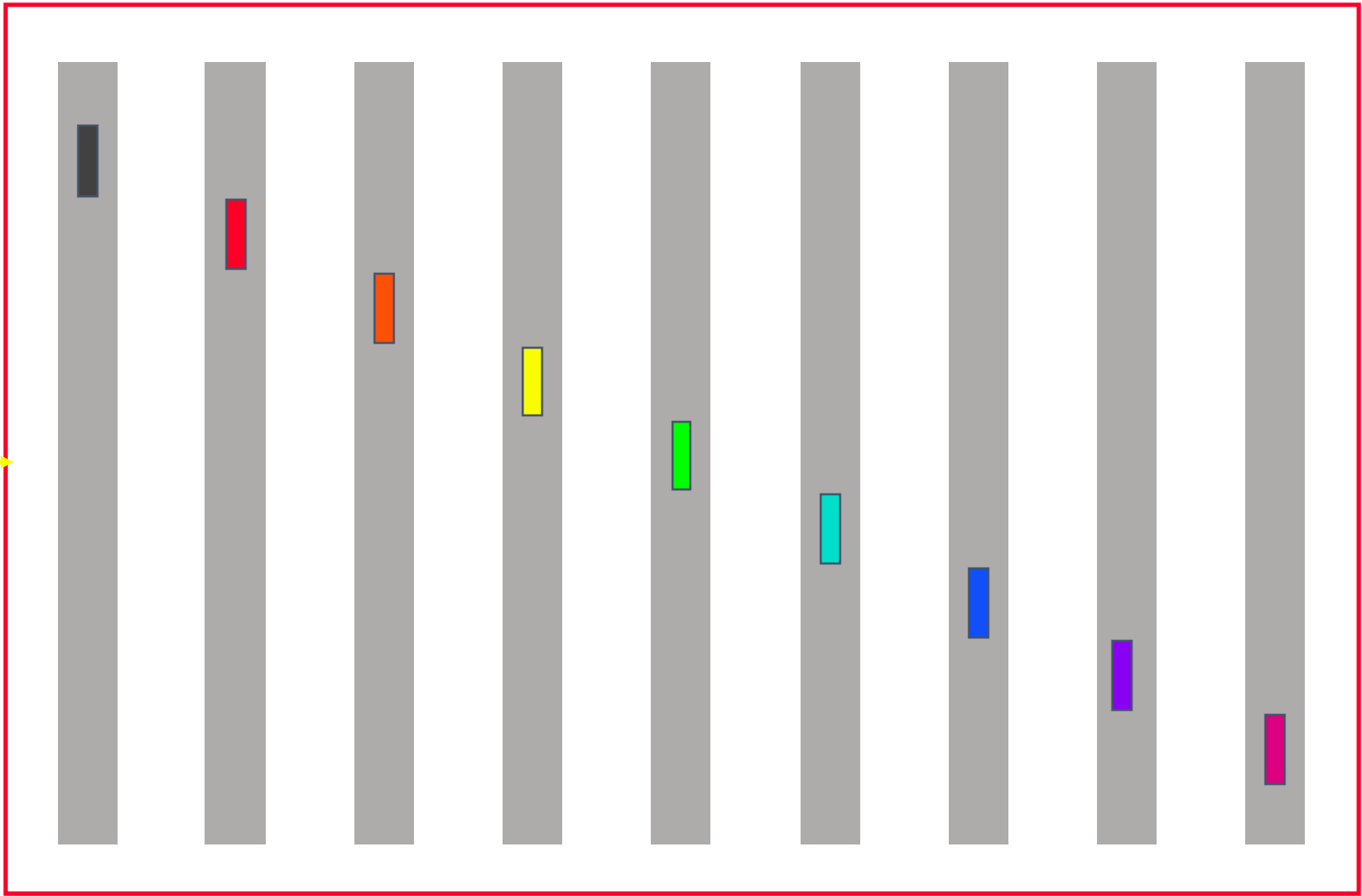
cost per steps

# Reduce-scatter

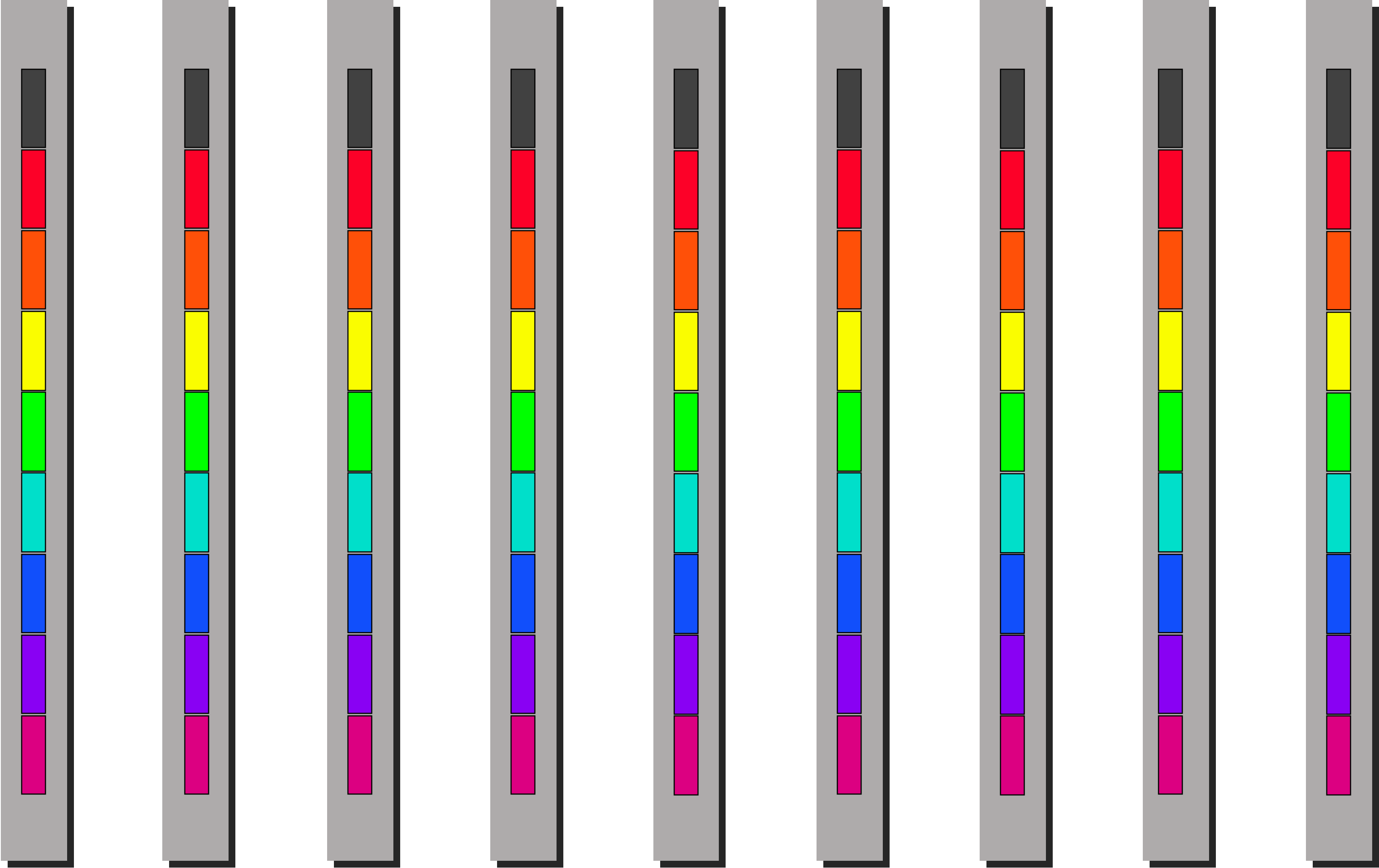
Before

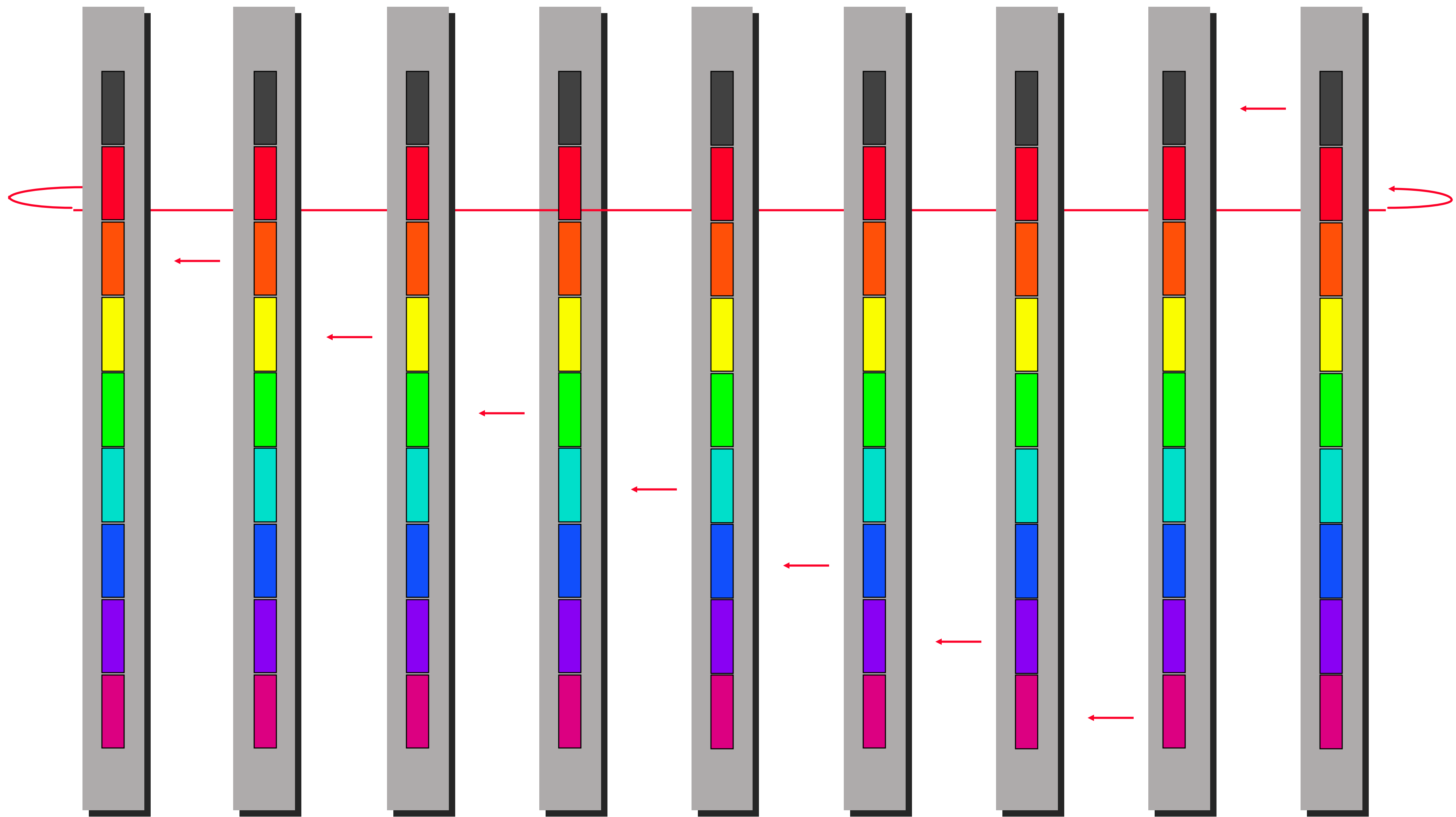


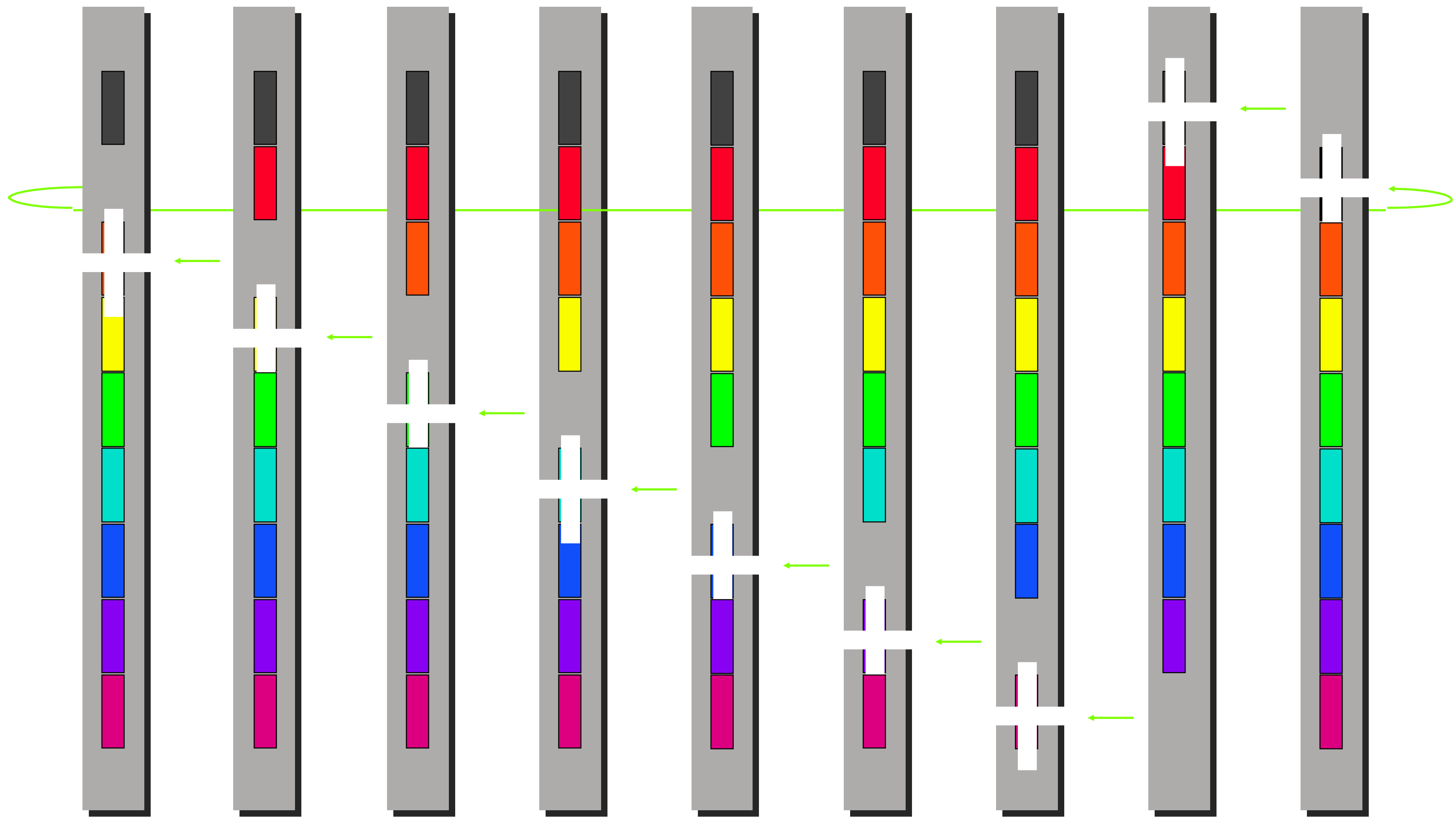
After

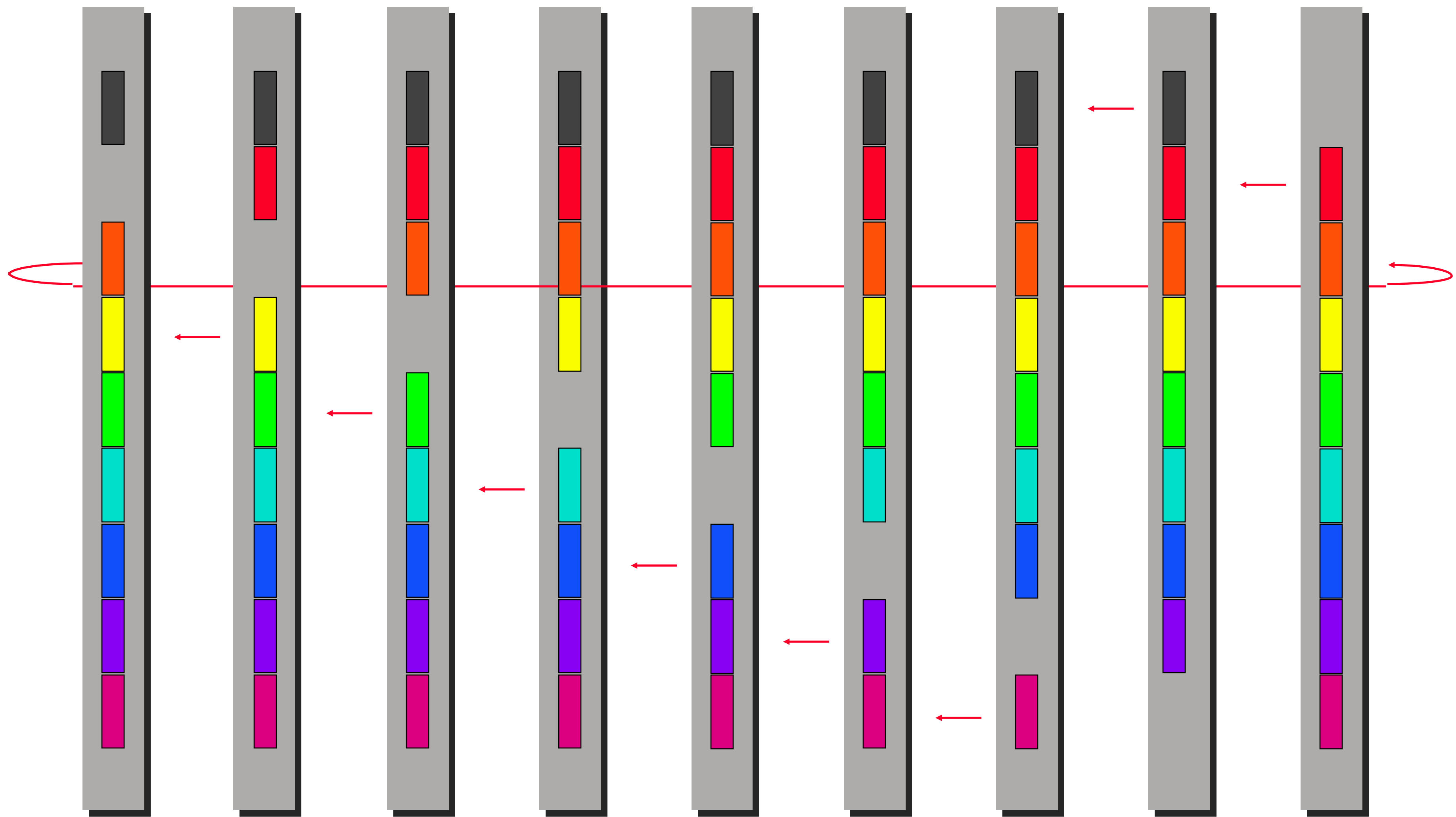


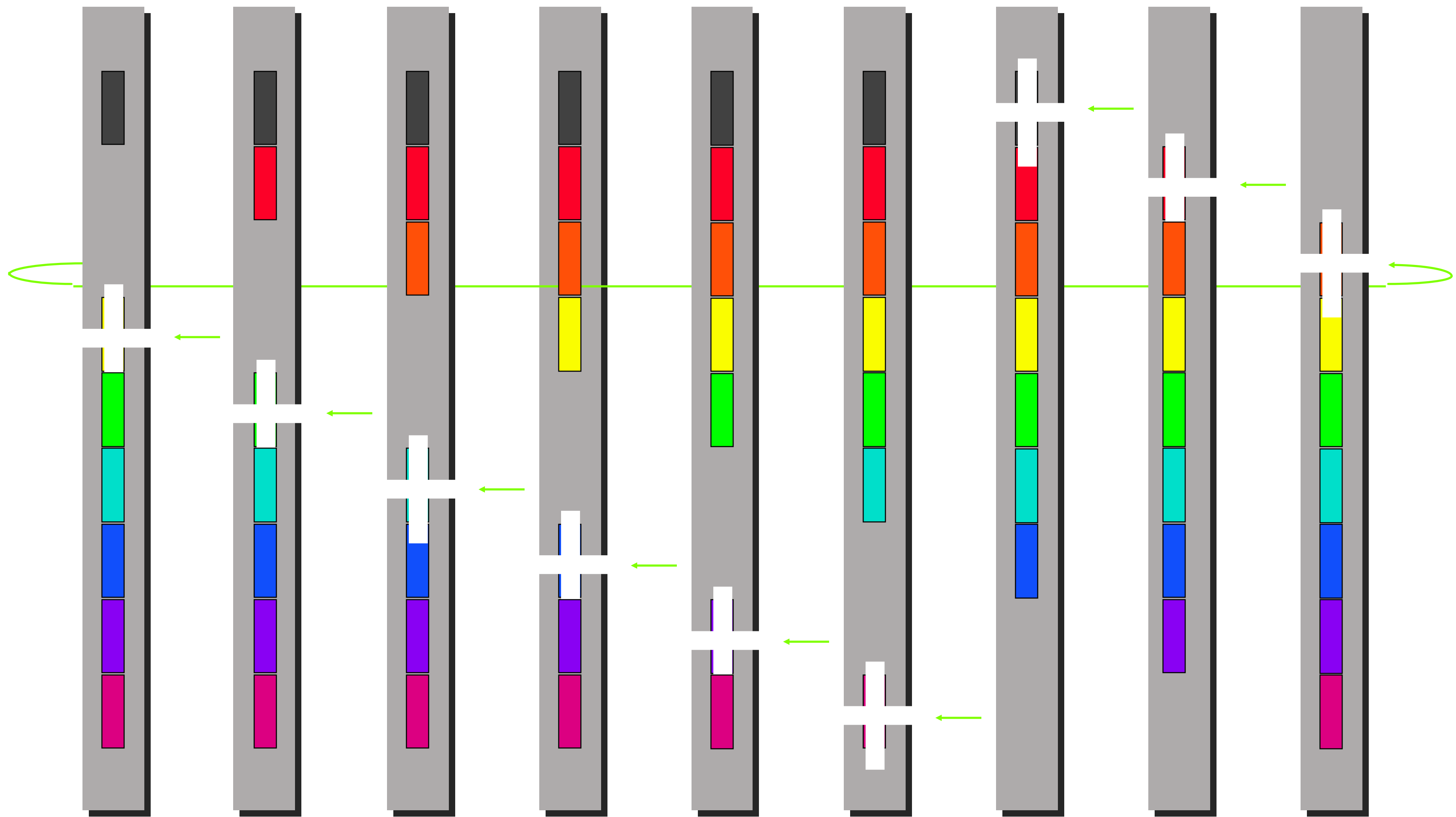


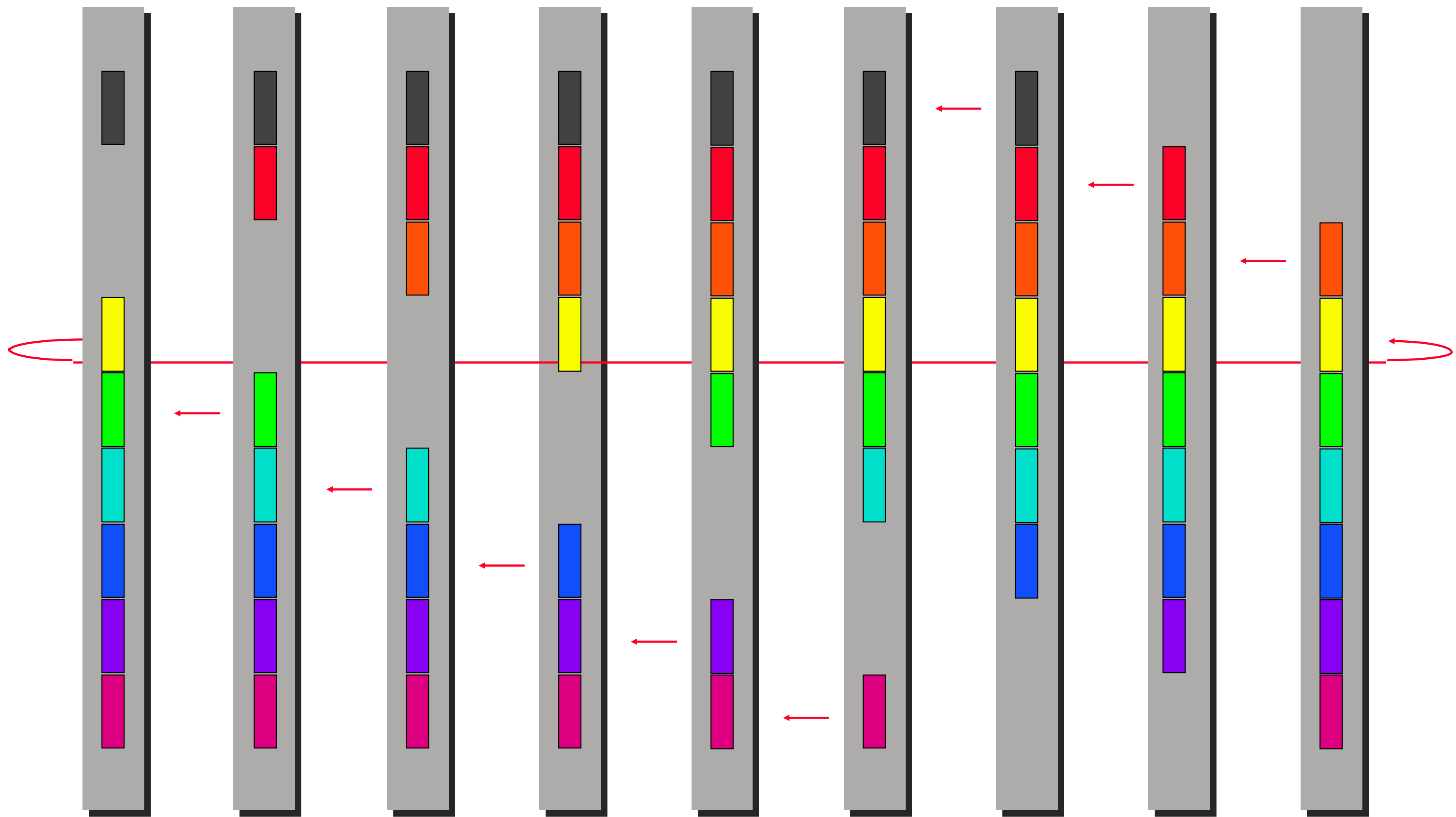


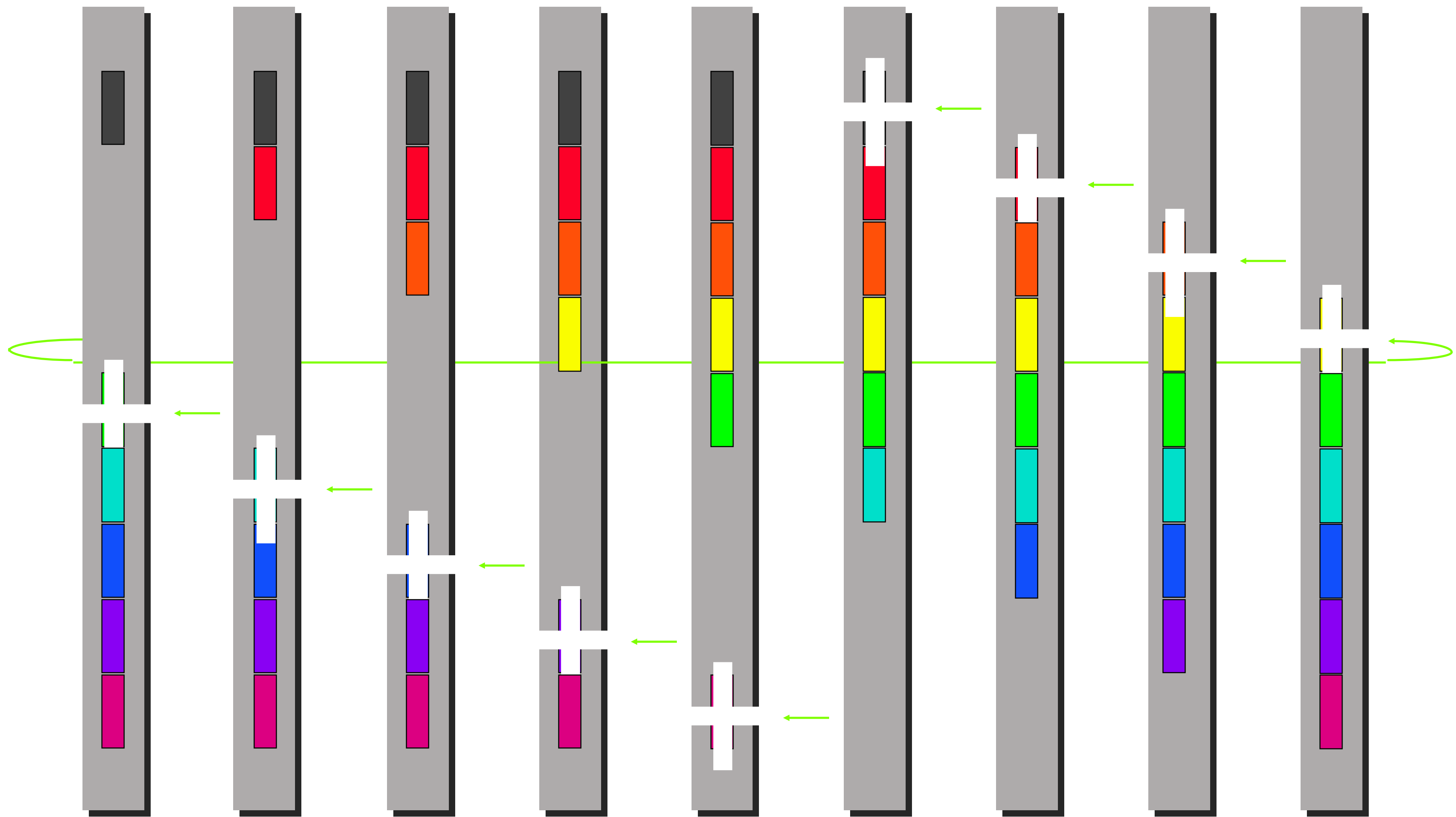


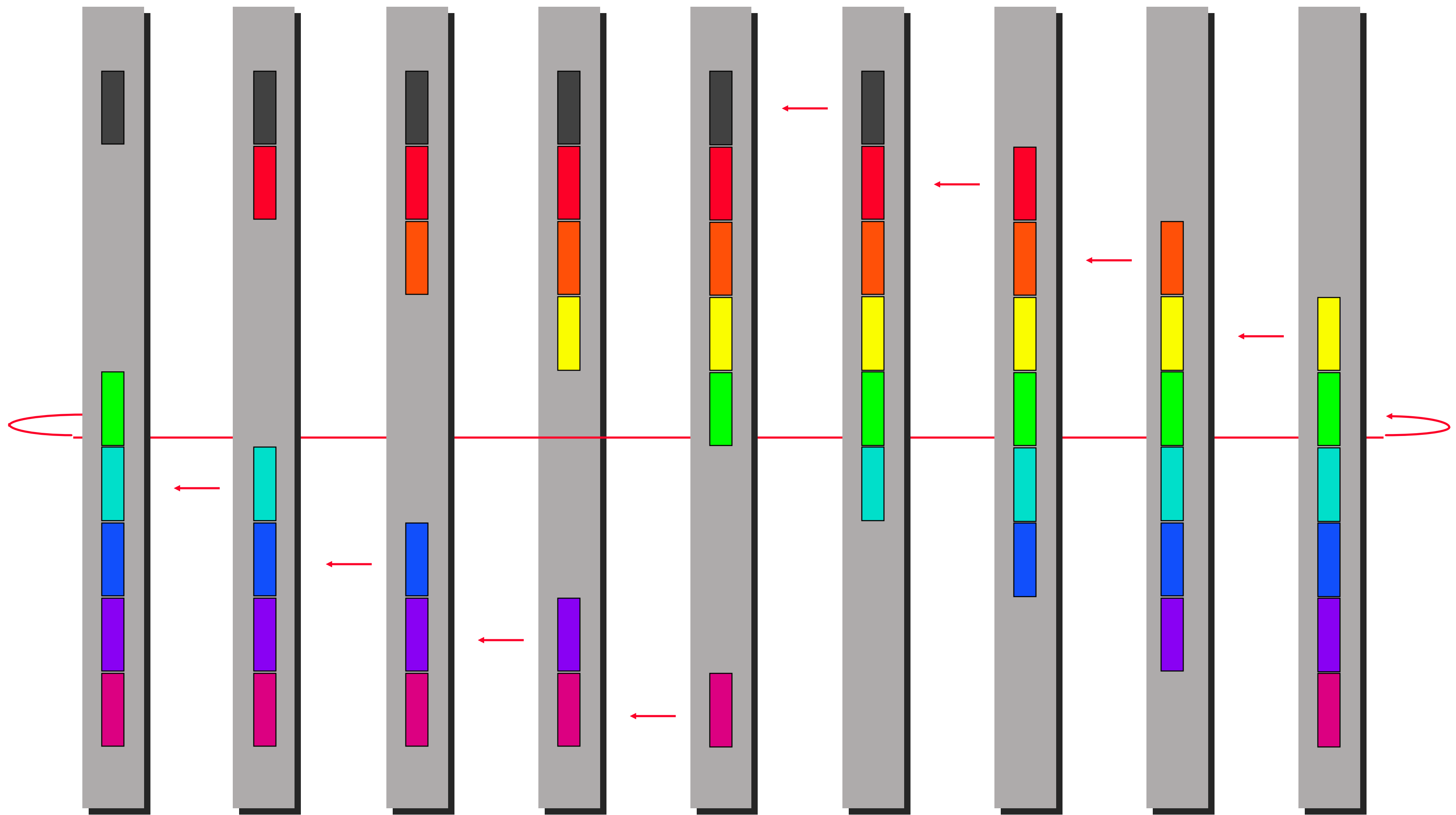




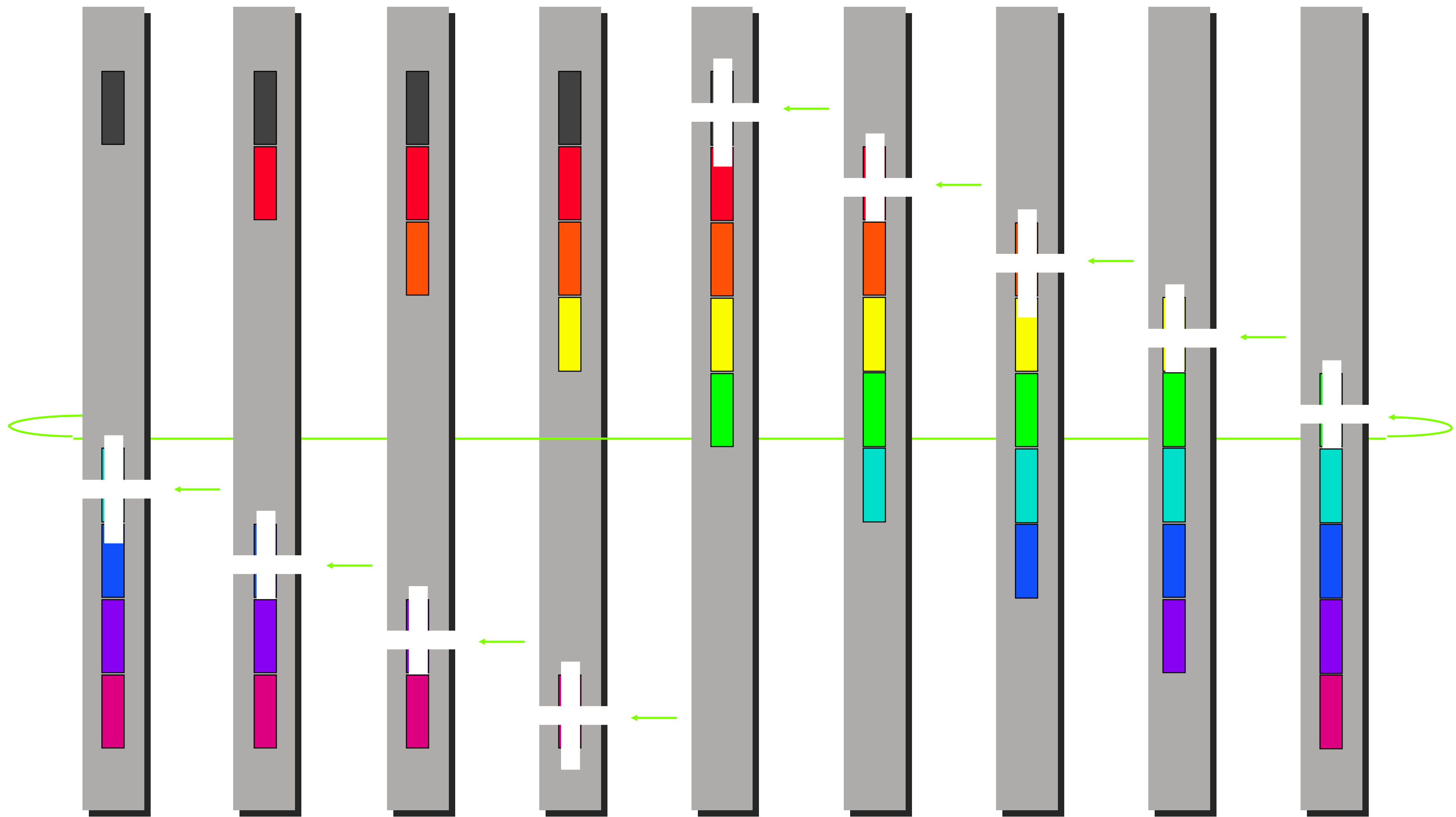


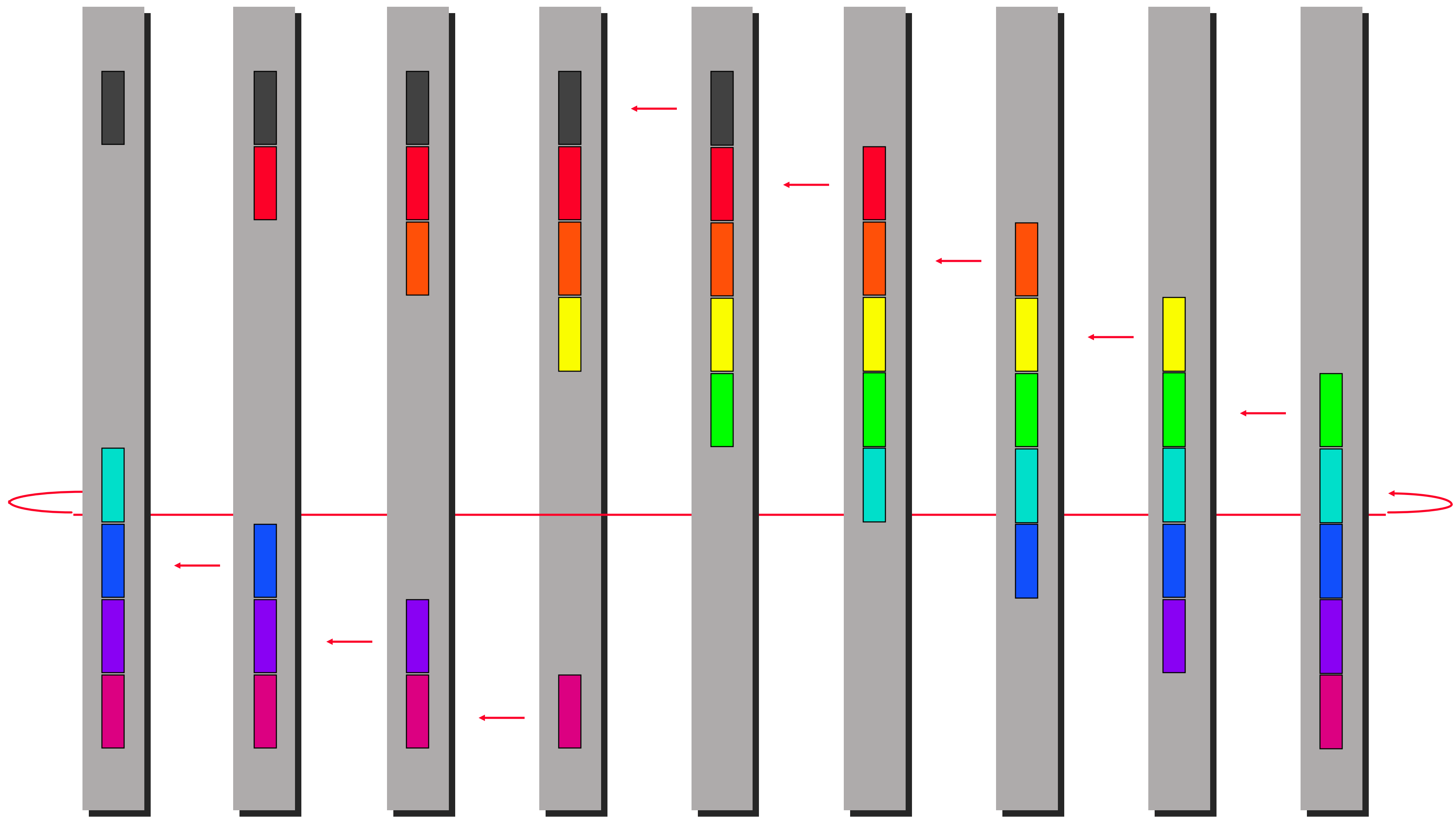


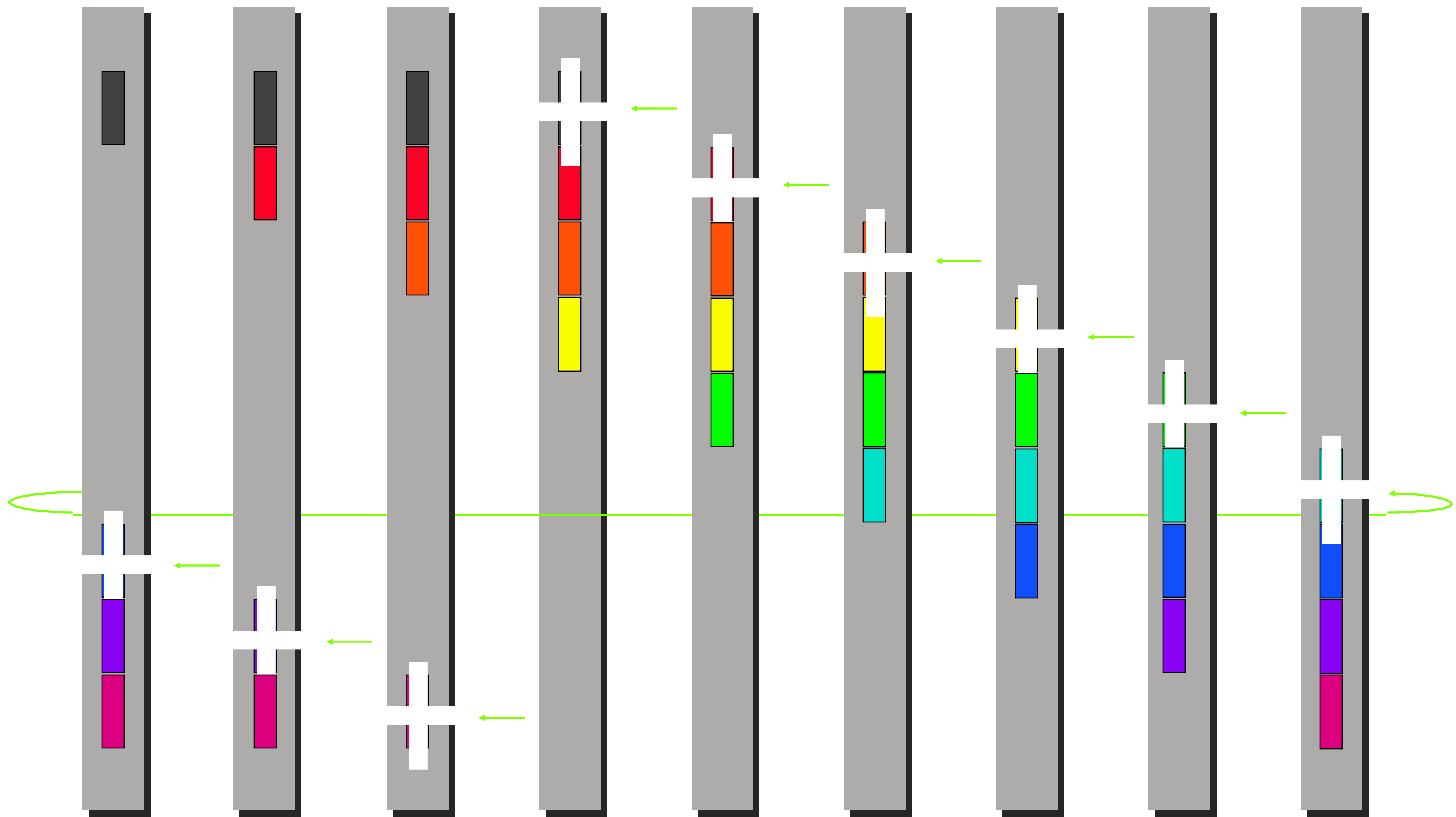


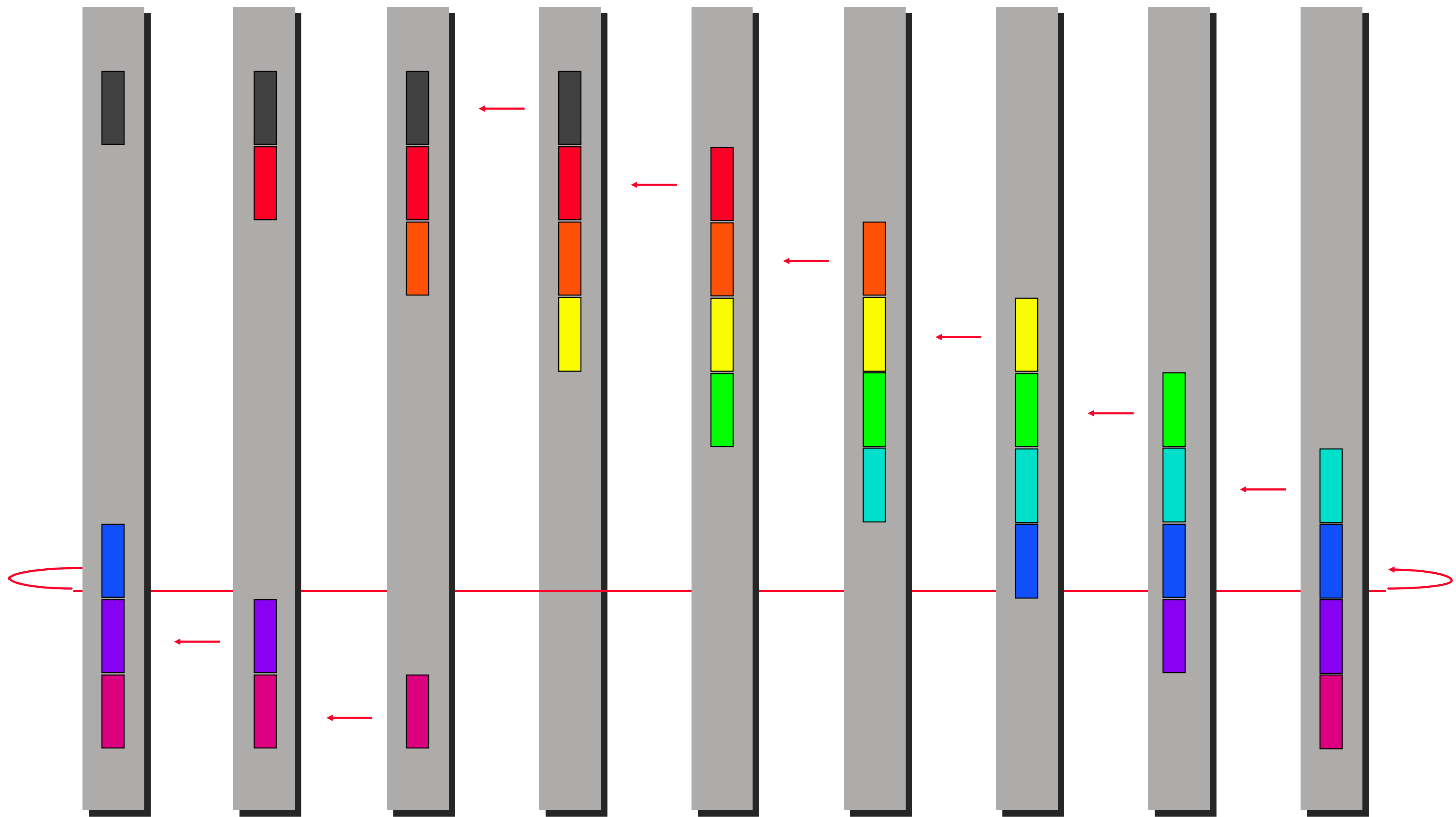


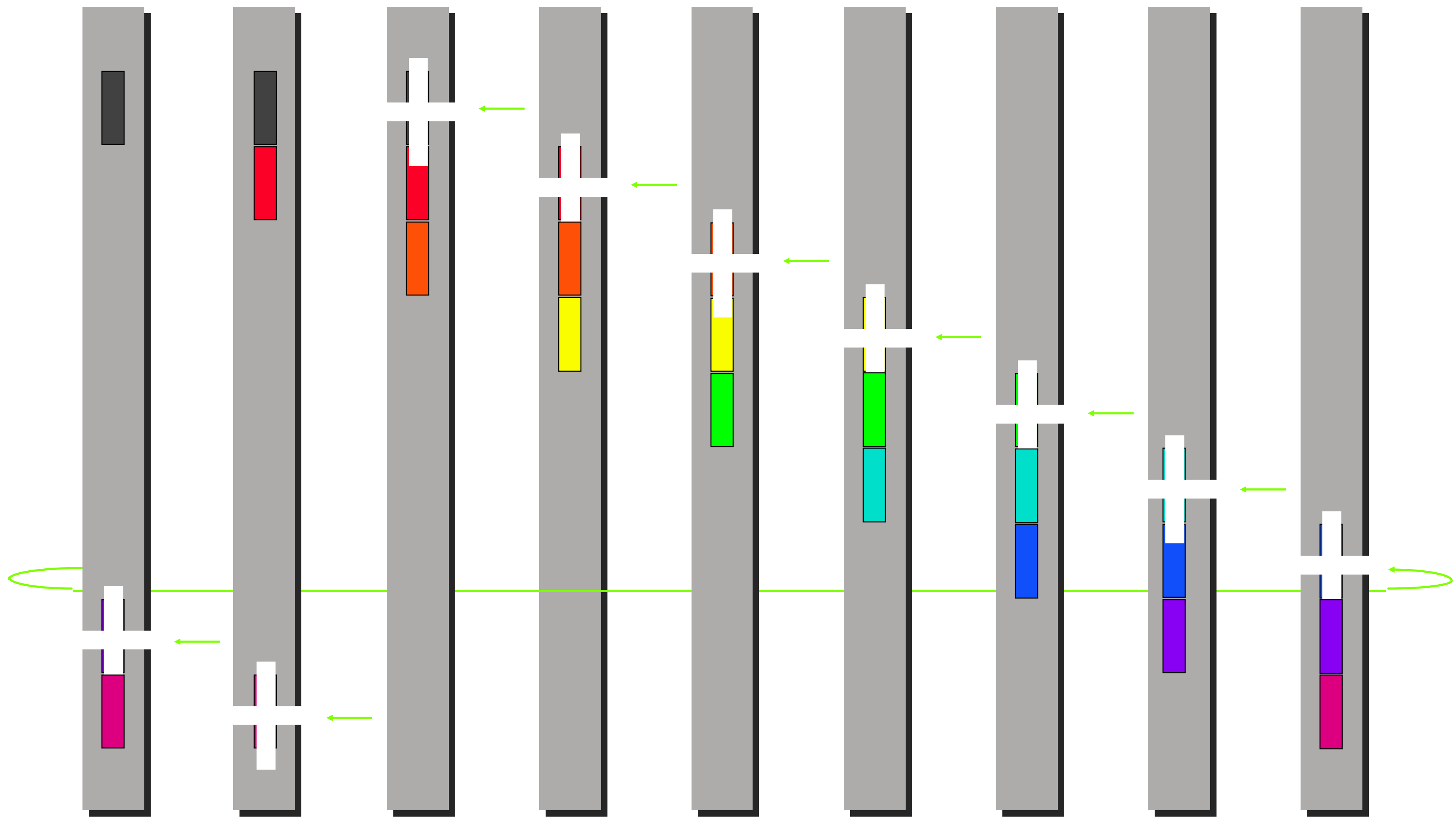


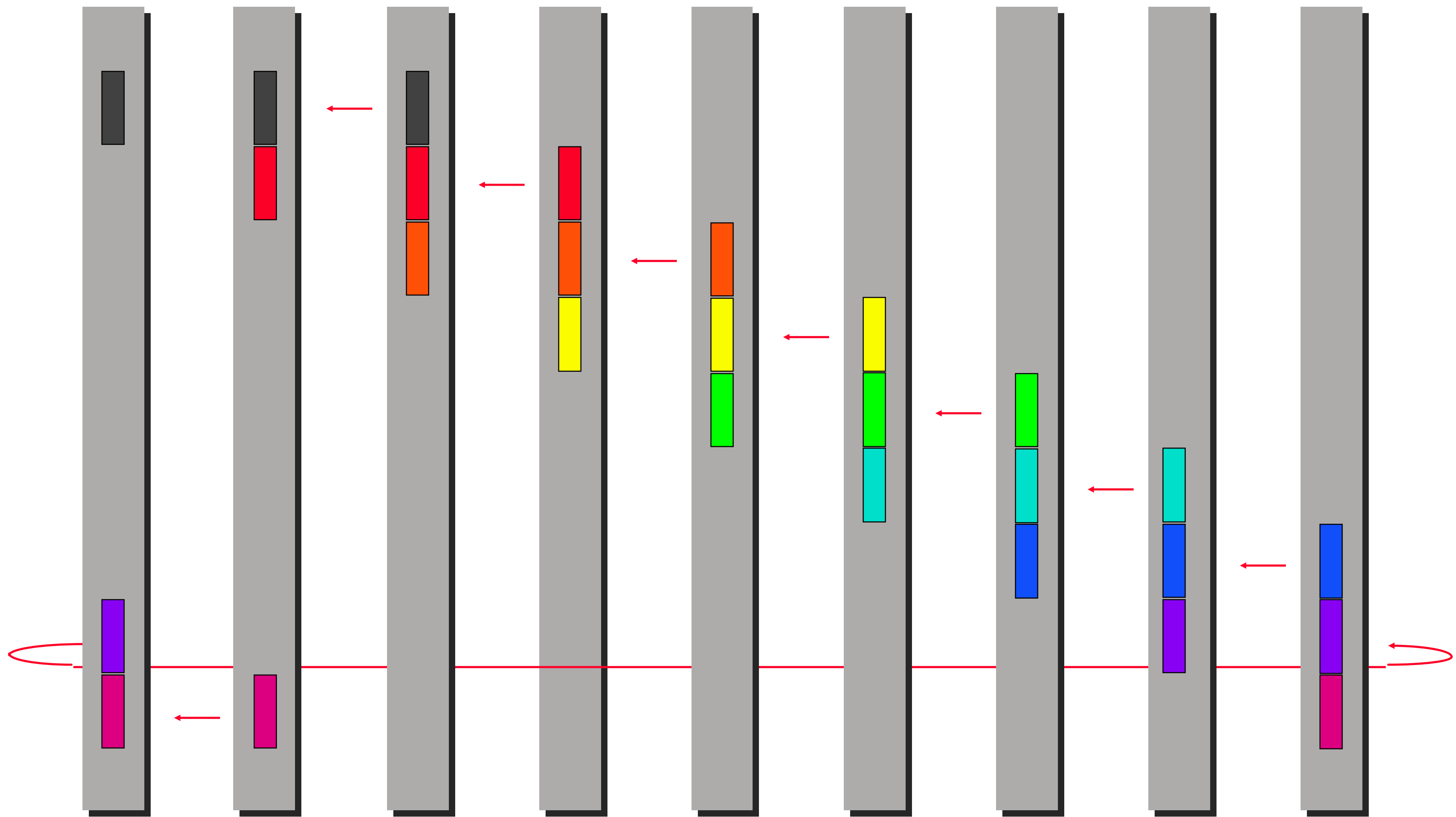


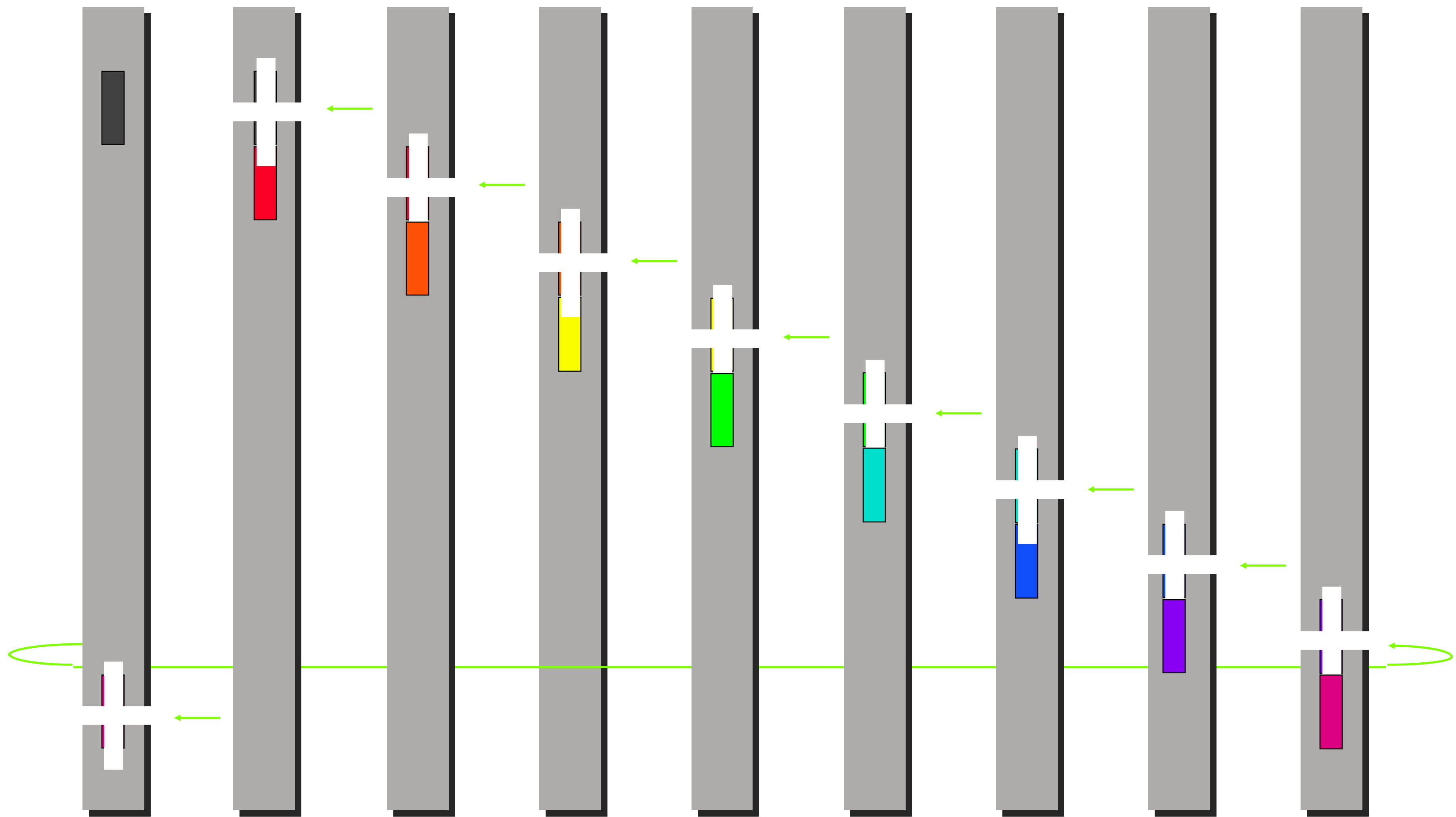


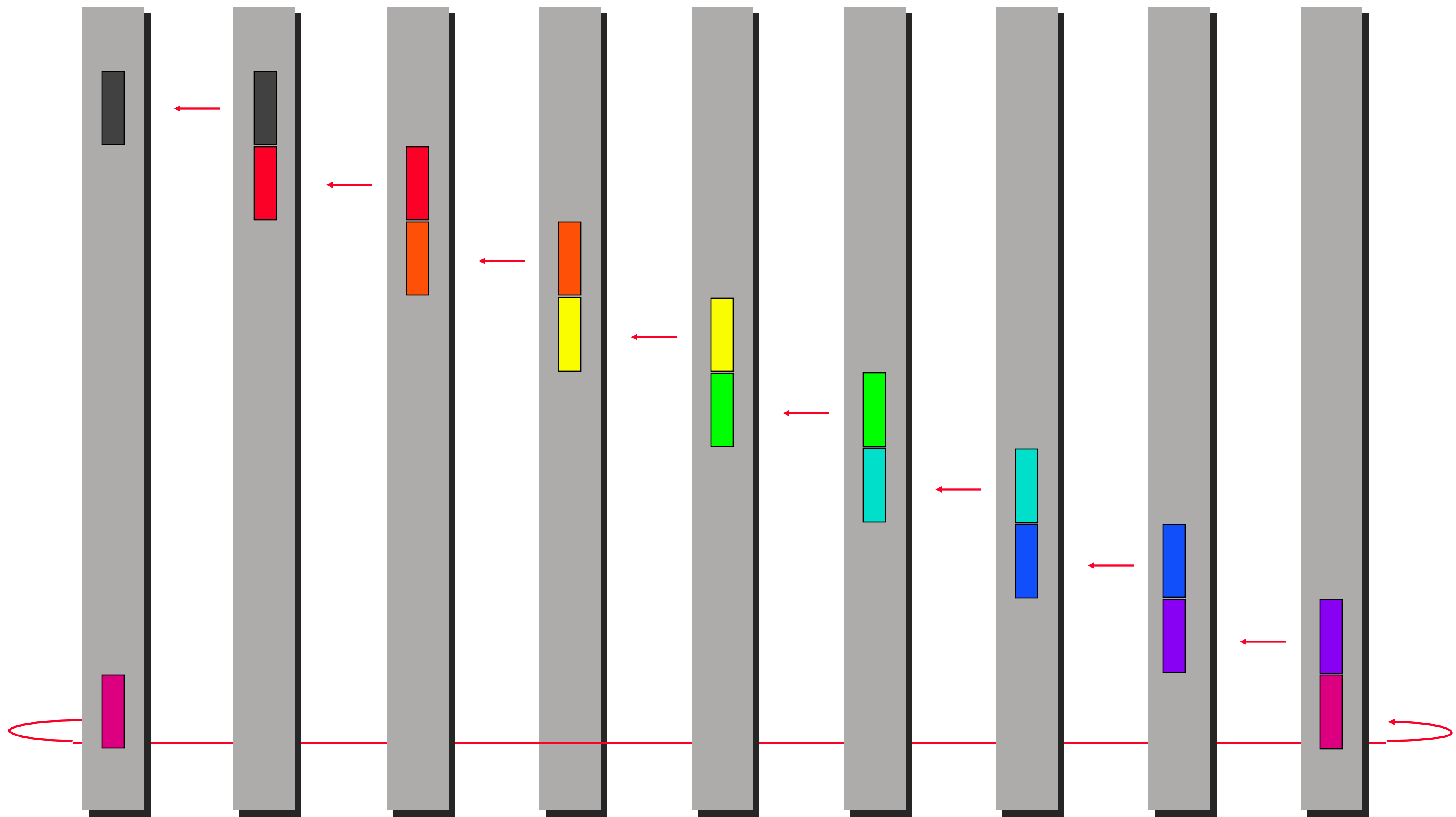




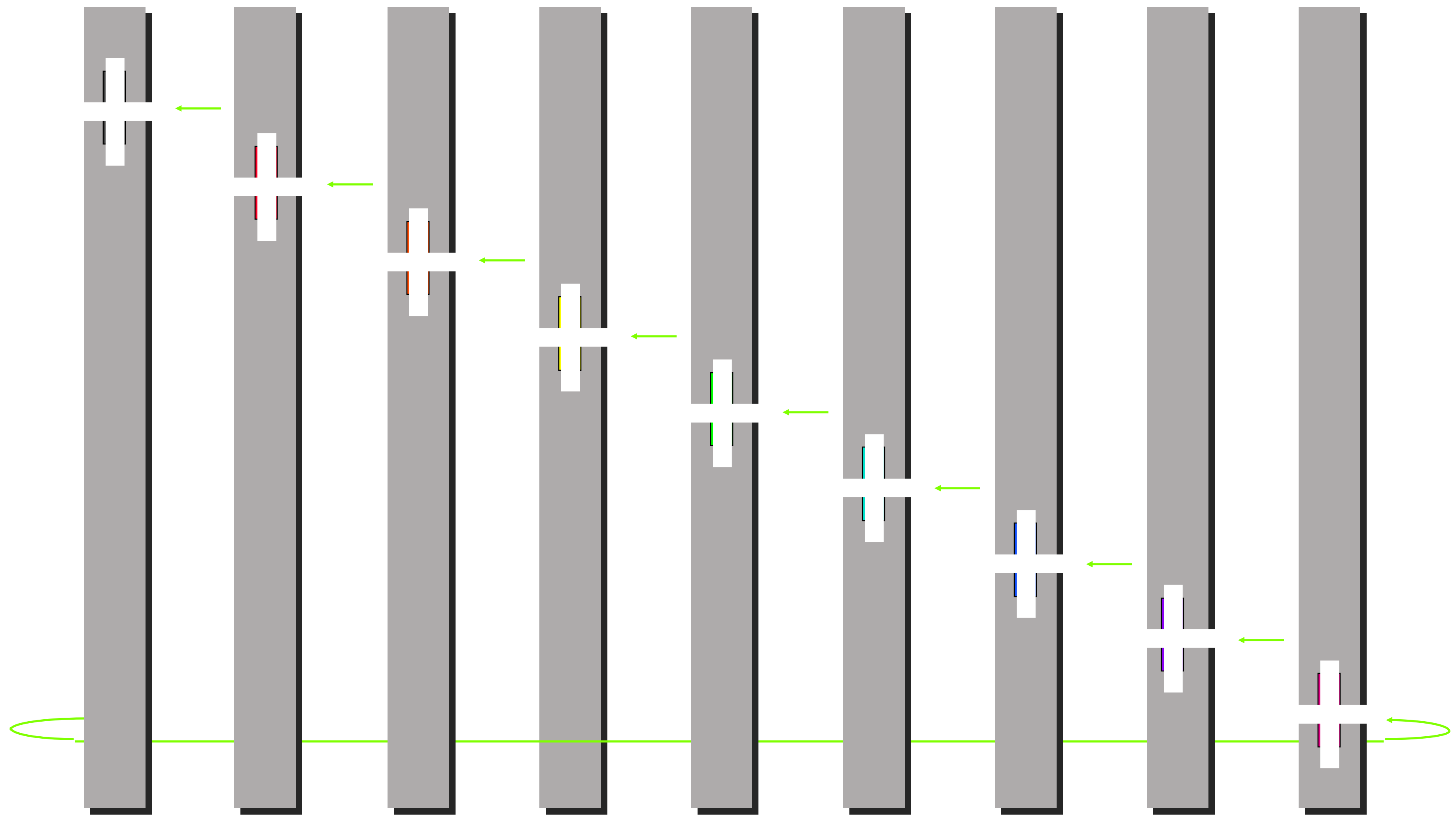


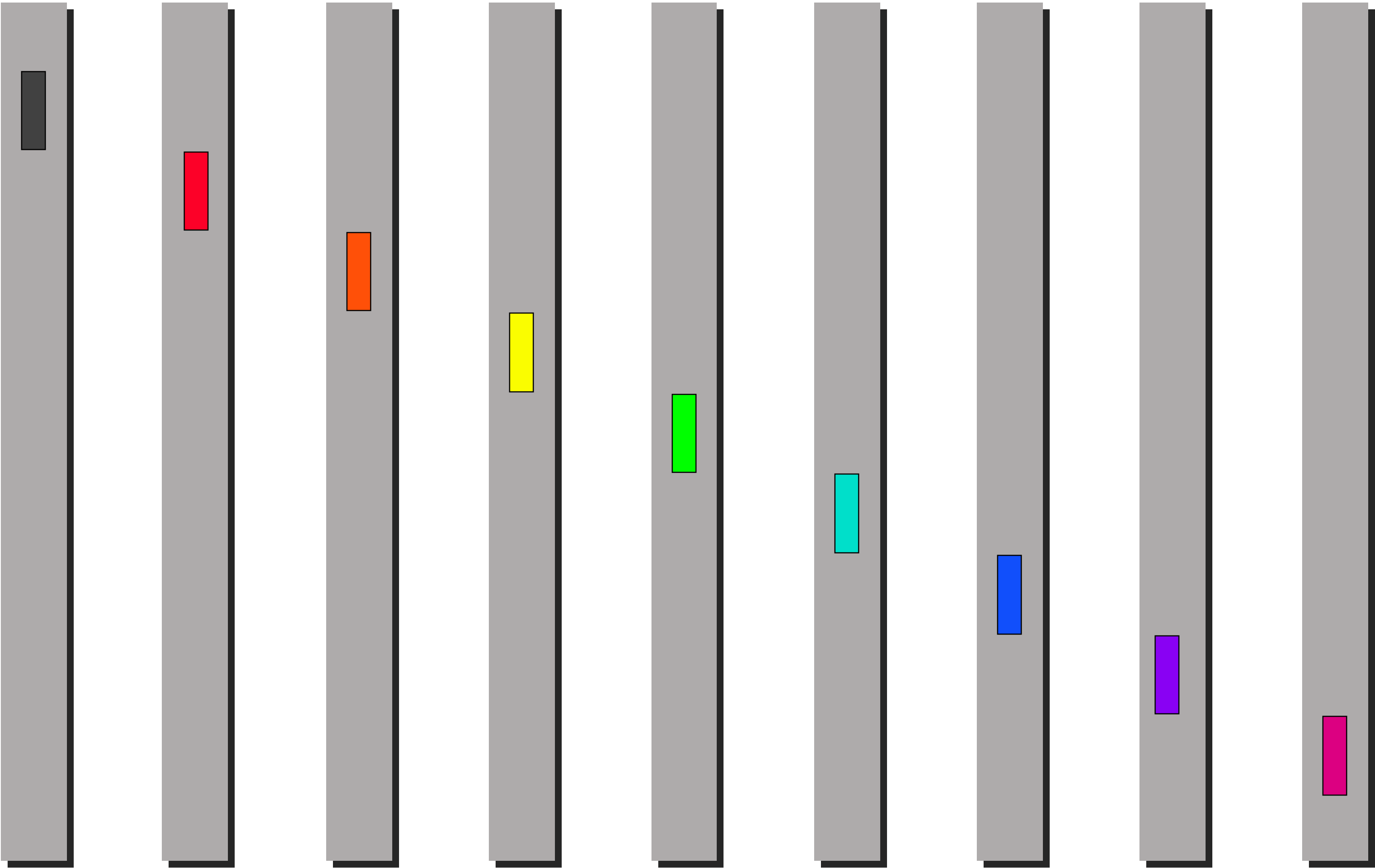












Cost

$$(p-1) \left( \alpha + \frac{n}{p} \beta + \frac{n}{p} \gamma \right) = (p-1)\alpha + \frac{p-1}{p} n \beta + \frac{p-1}{p} n \gamma$$

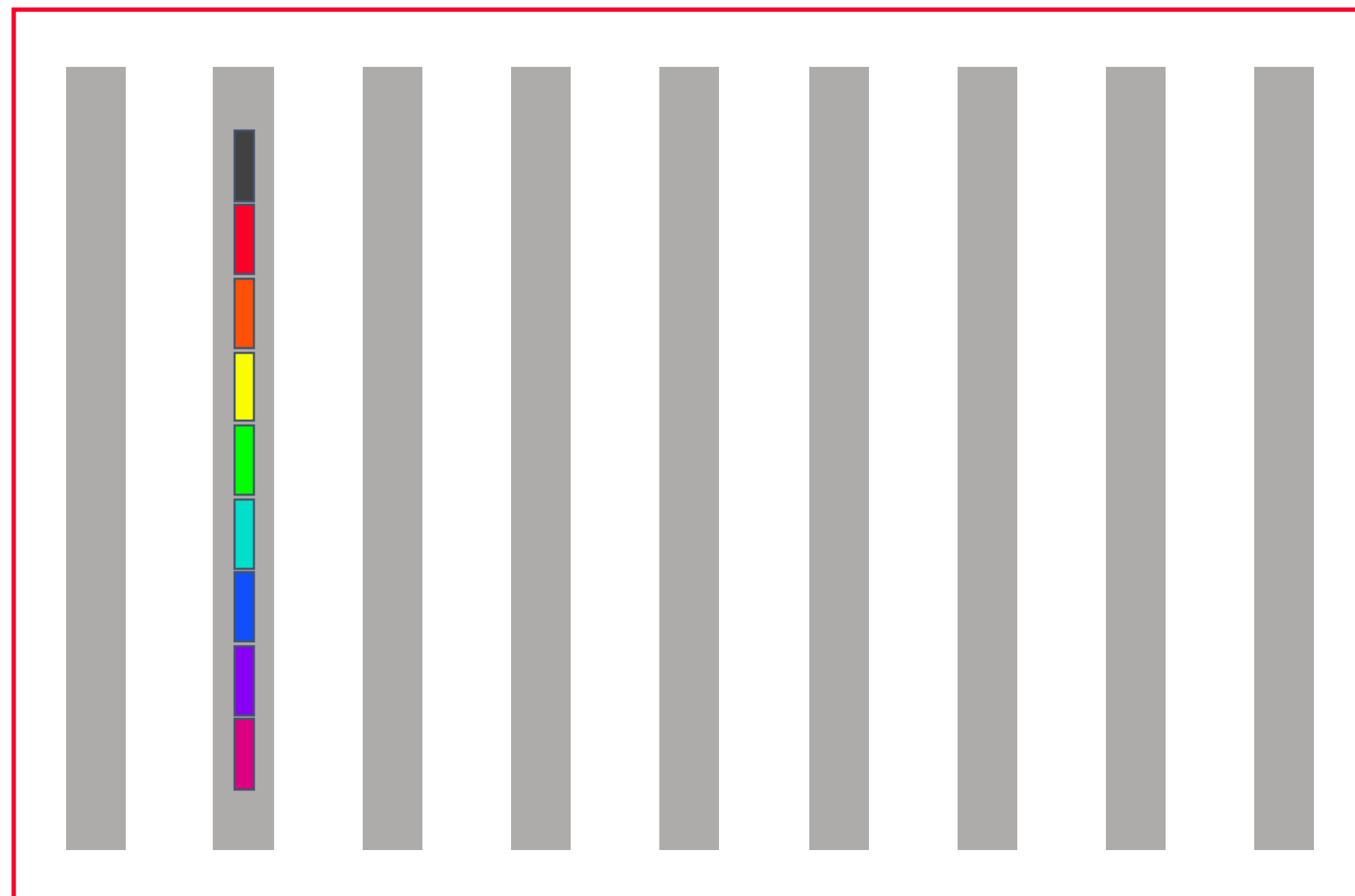
number of steps

cost per steps

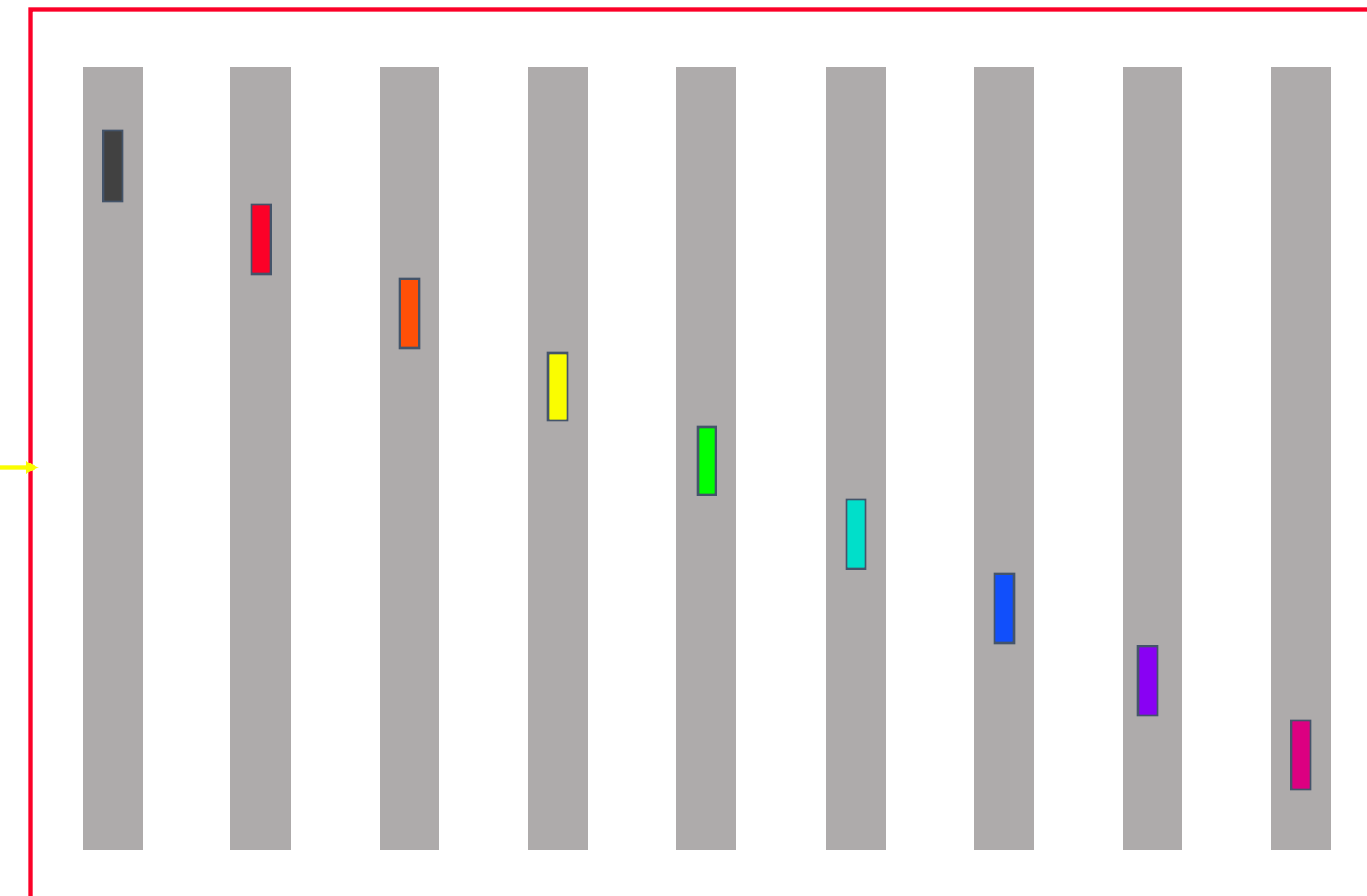
# Scatter: Can Ring Be Better?

Notice: Scatter as implemented before using *MST* was optimal in Bandwidth as well (How to Prove?)

Before



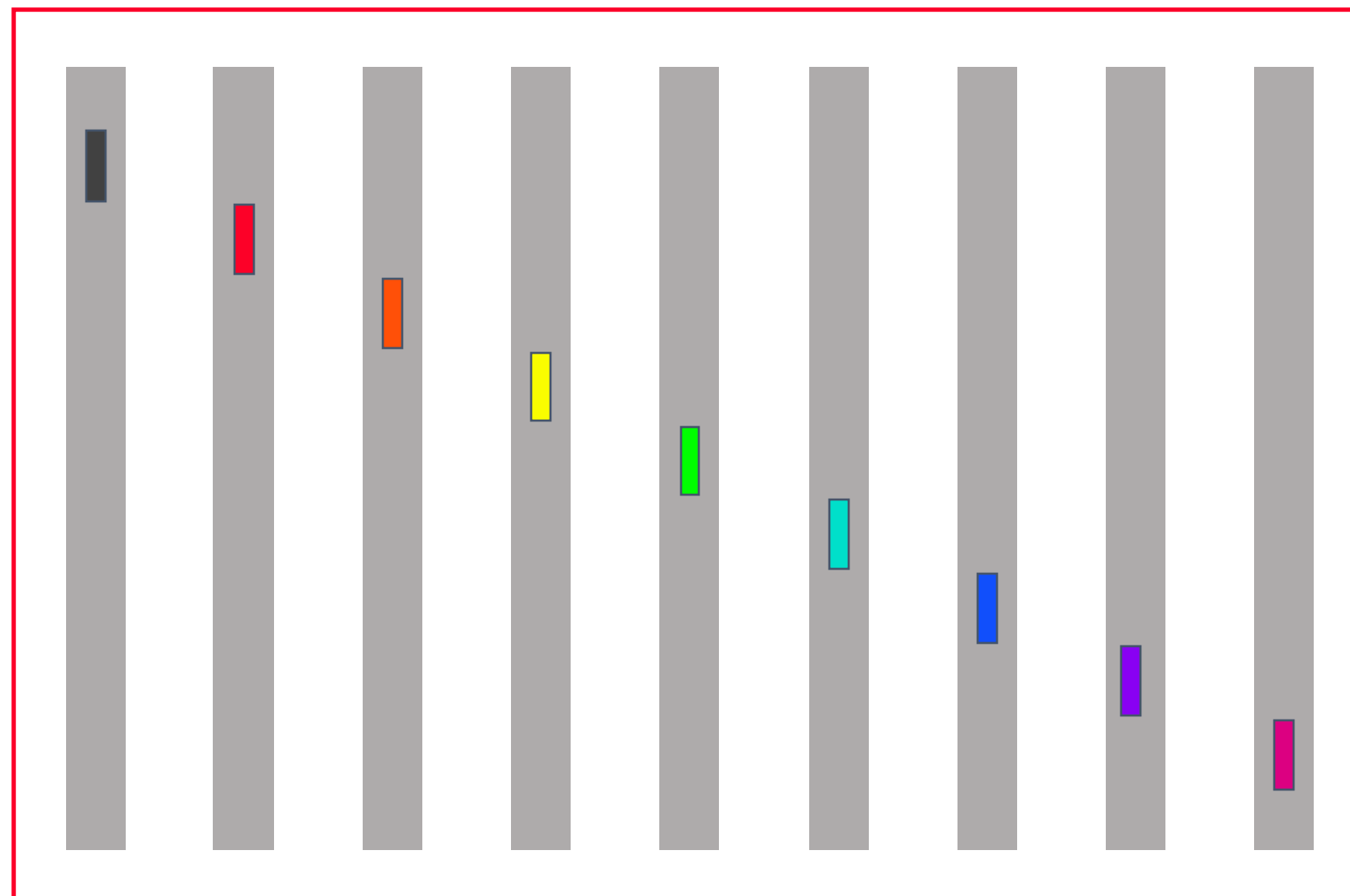
After



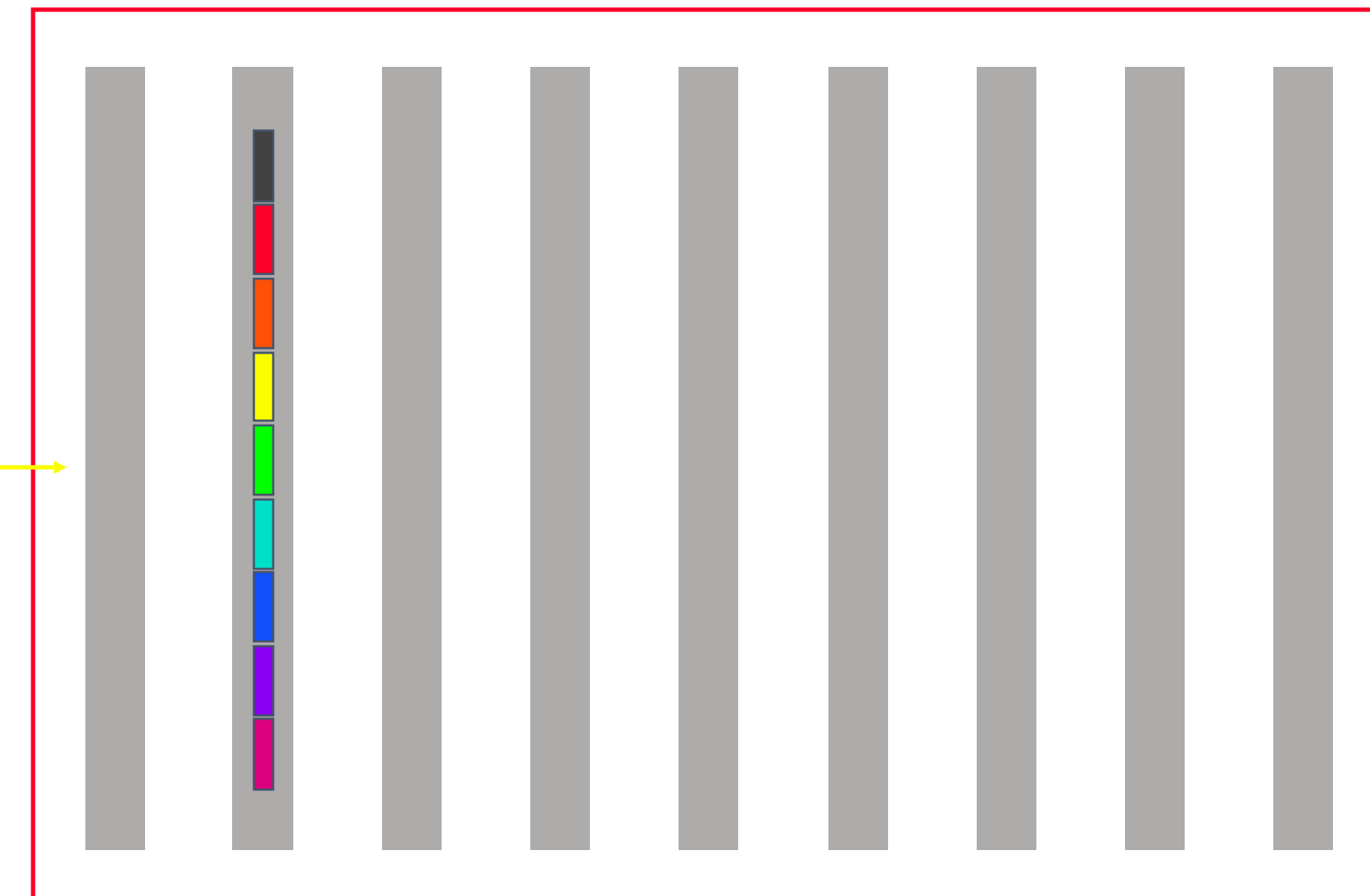
# Gather

Notice: Gather as implemented before using MST was optimal in bandwidth as well (how to prove?)

Before

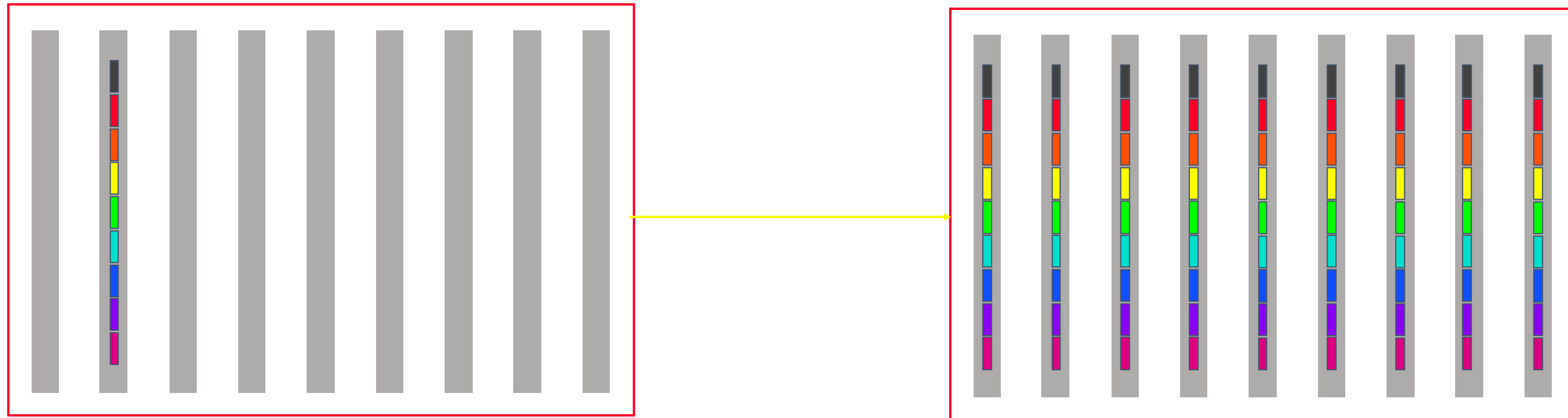


After

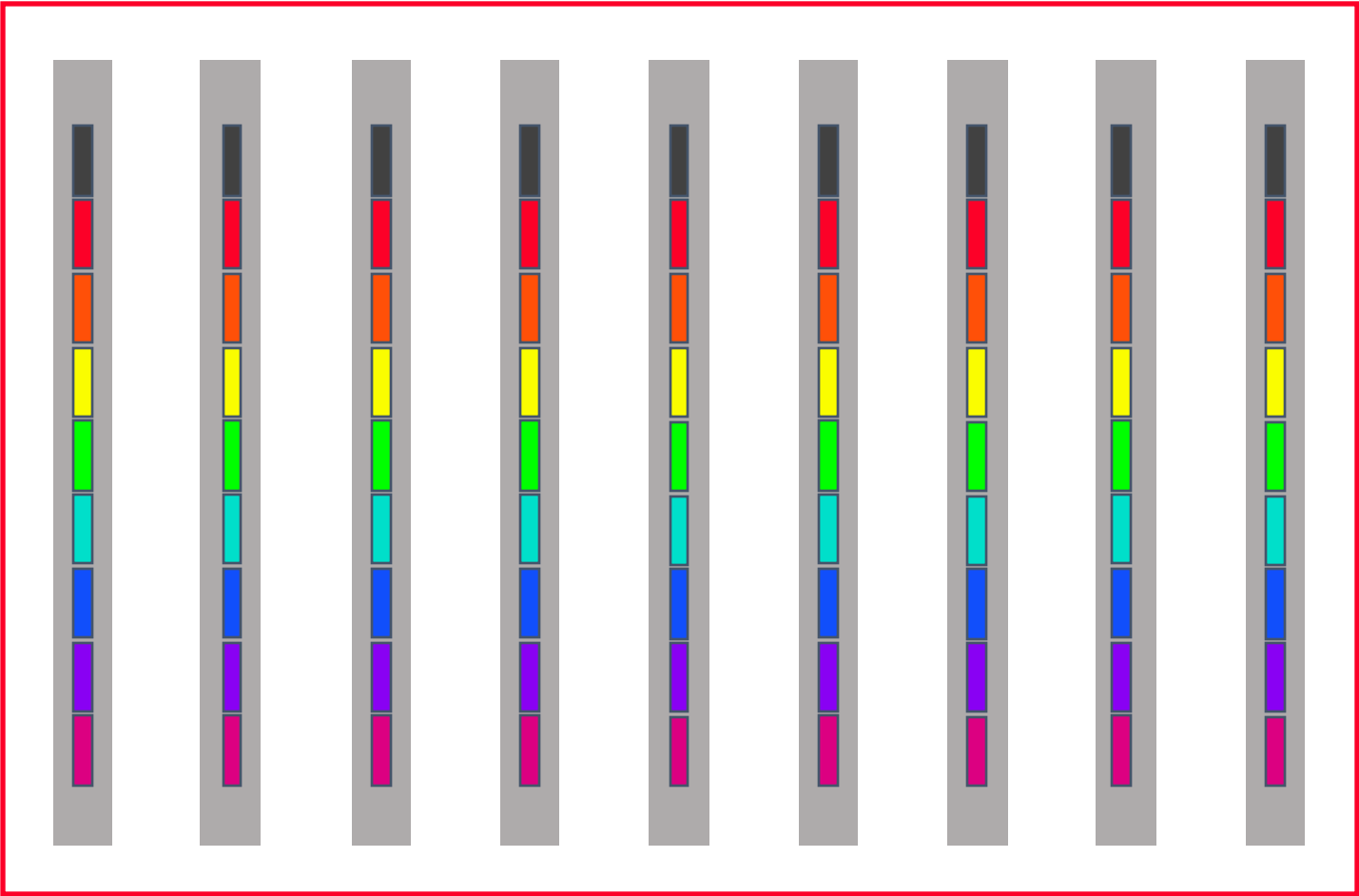


# Using the building blocks

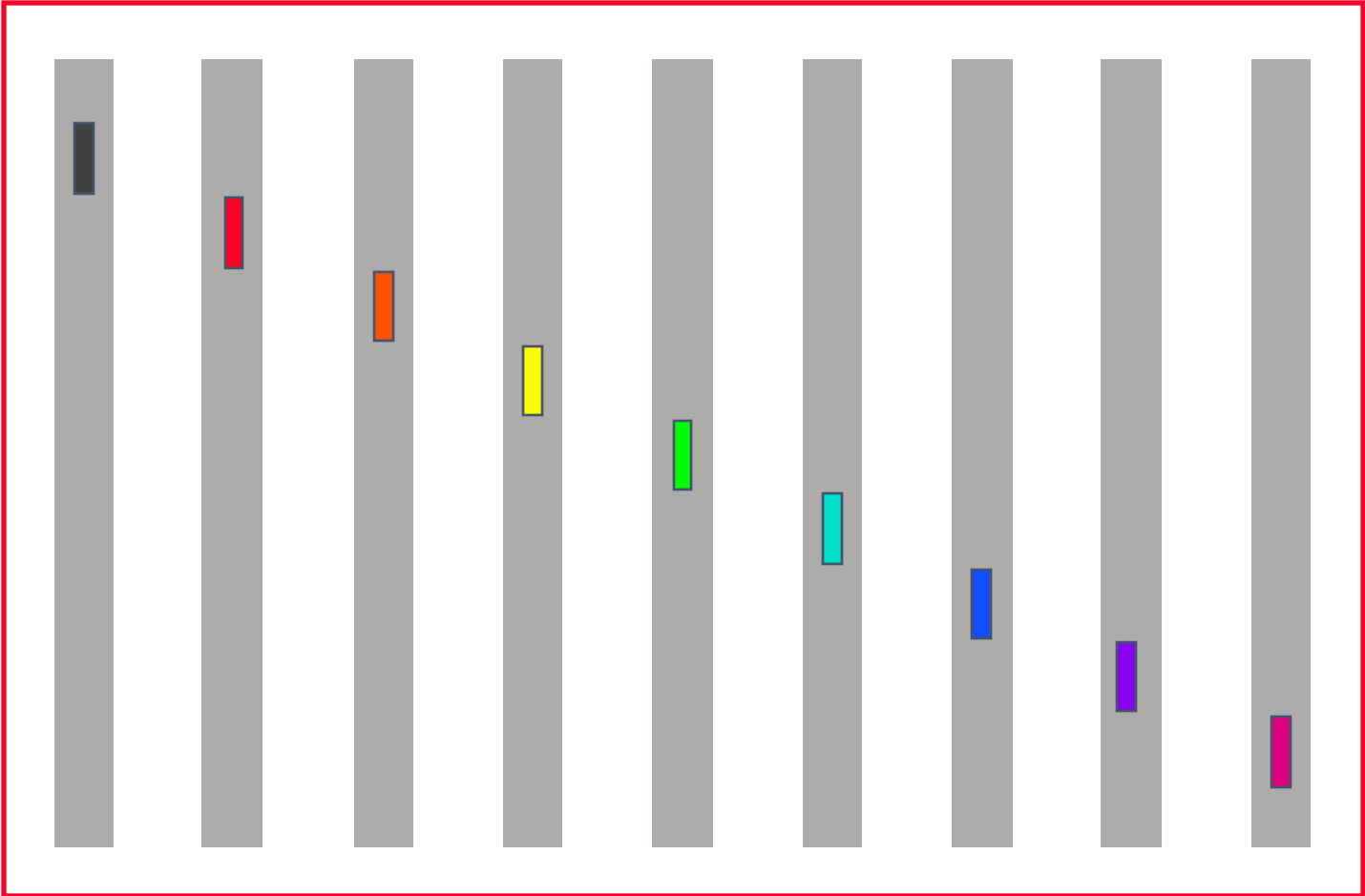
# Broadcast (Large Message)



# Broadcast (Large Message)

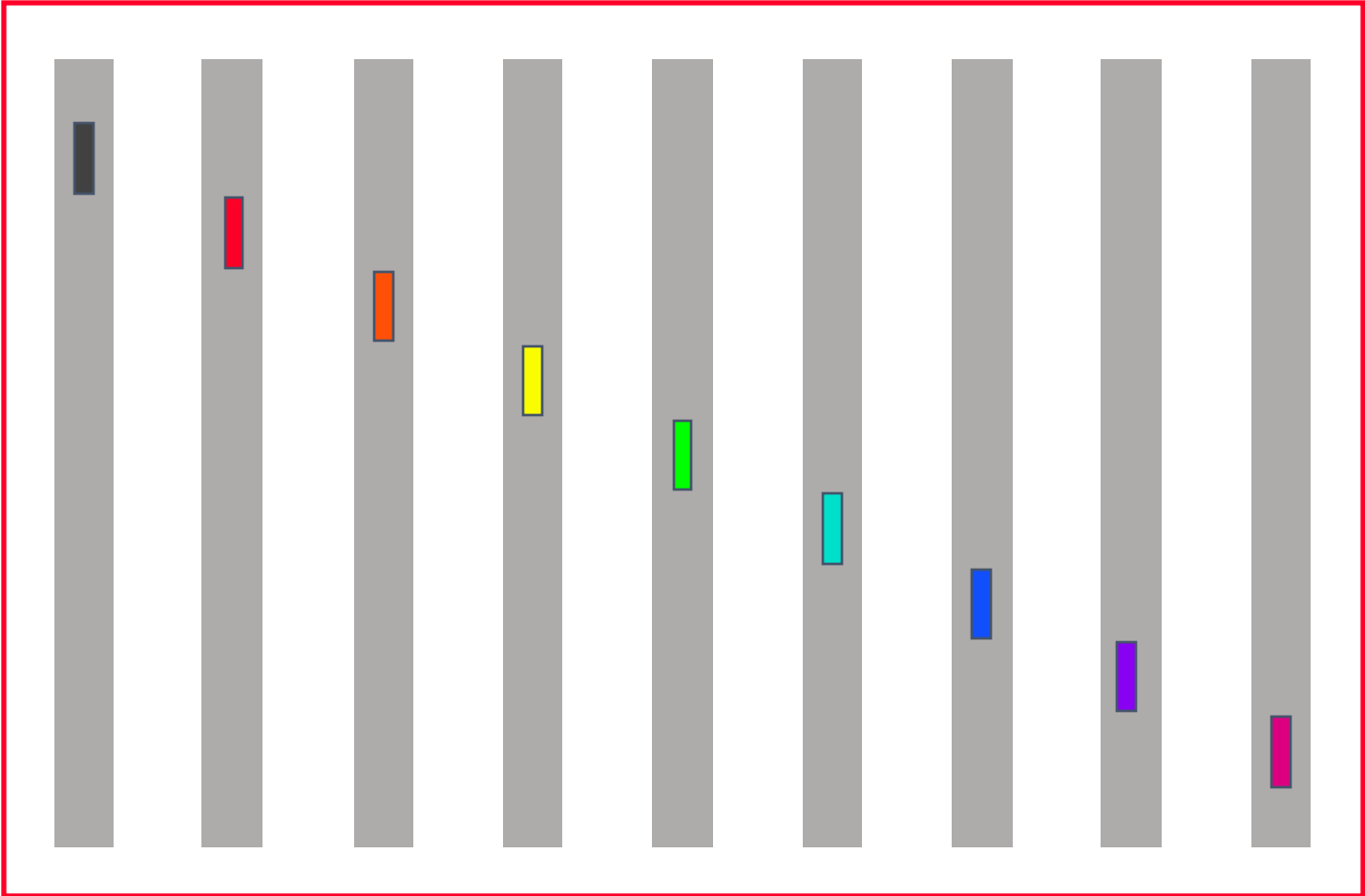
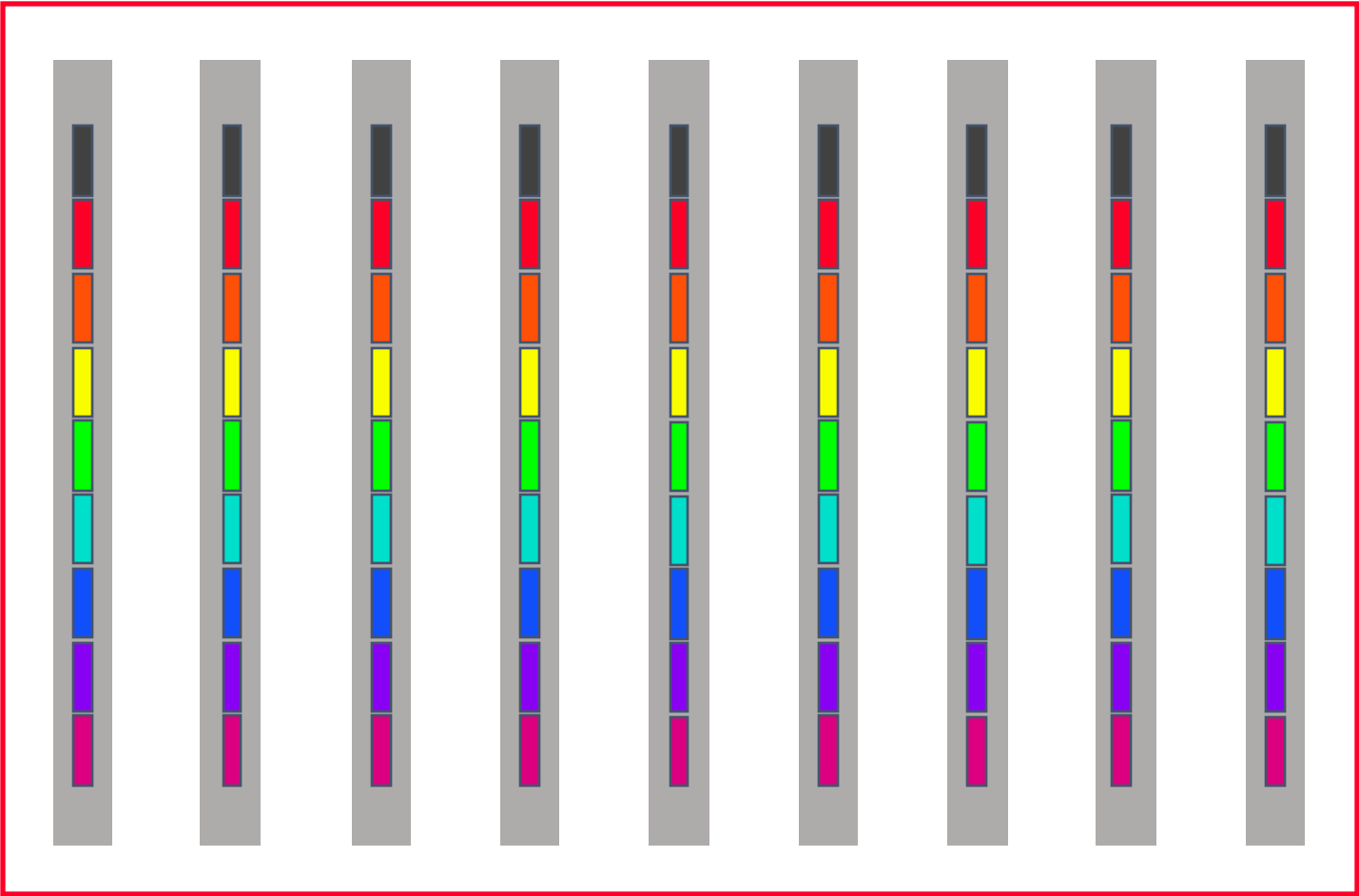


Scatter





# Broadcast (Large Message)



Allgather

# Cost of scatter/allgather broadcast

- Assumption: power of two number of nodes

scatter  $\log(p)\alpha + \frac{p-1}{p}n\beta$

allgather  $(p-1)\alpha + \frac{p-1}{p}n\beta$

---

$$(\log(p) + p - 1)\alpha + 2\frac{p-1}{p}n\beta$$

# Cost of scatter/allgather broadcast

- Assumption: power of two number of nodes

scatter  $\log(p)\alpha + \frac{p-1}{p}n\beta$

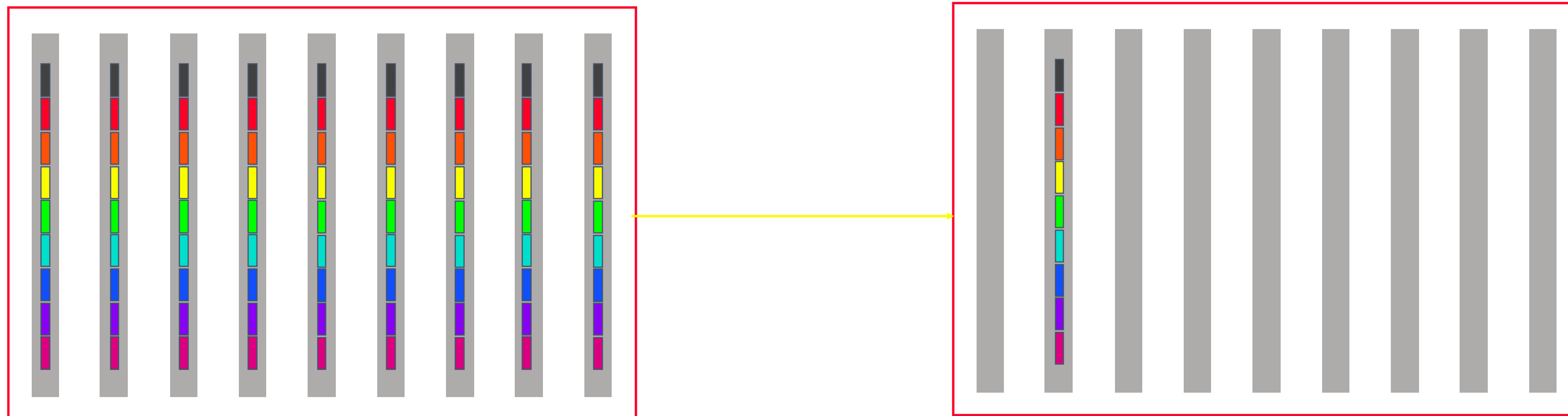
allgather  $(p-1)\alpha + \frac{p-1}{p}n\beta$

---

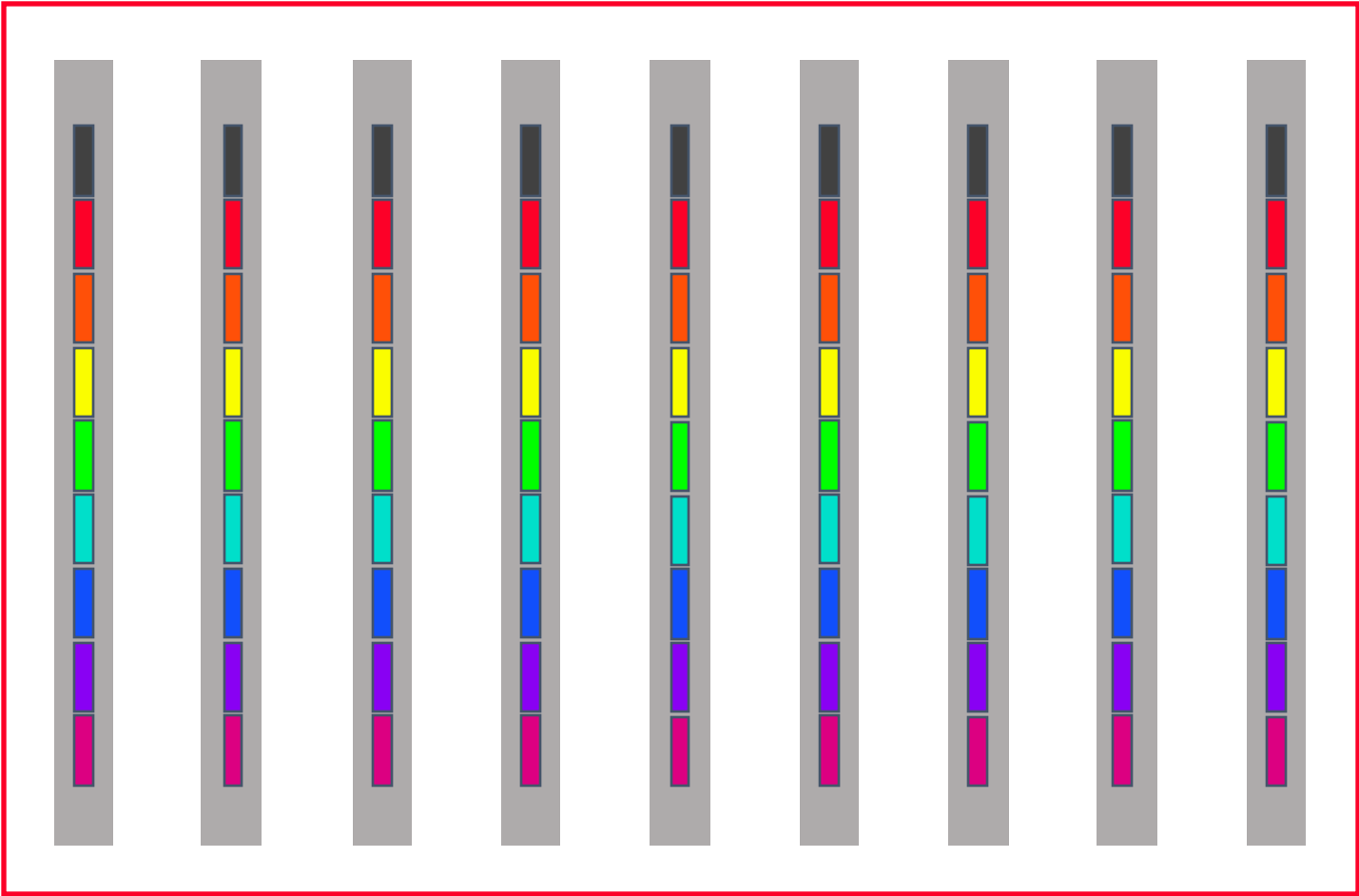

$$(\log(p) + p - 1)\alpha + 2\frac{p-1}{p}n\beta$$

Vs. MST broadcast:  $\lceil \log(p) \rceil (\alpha + n\beta)$

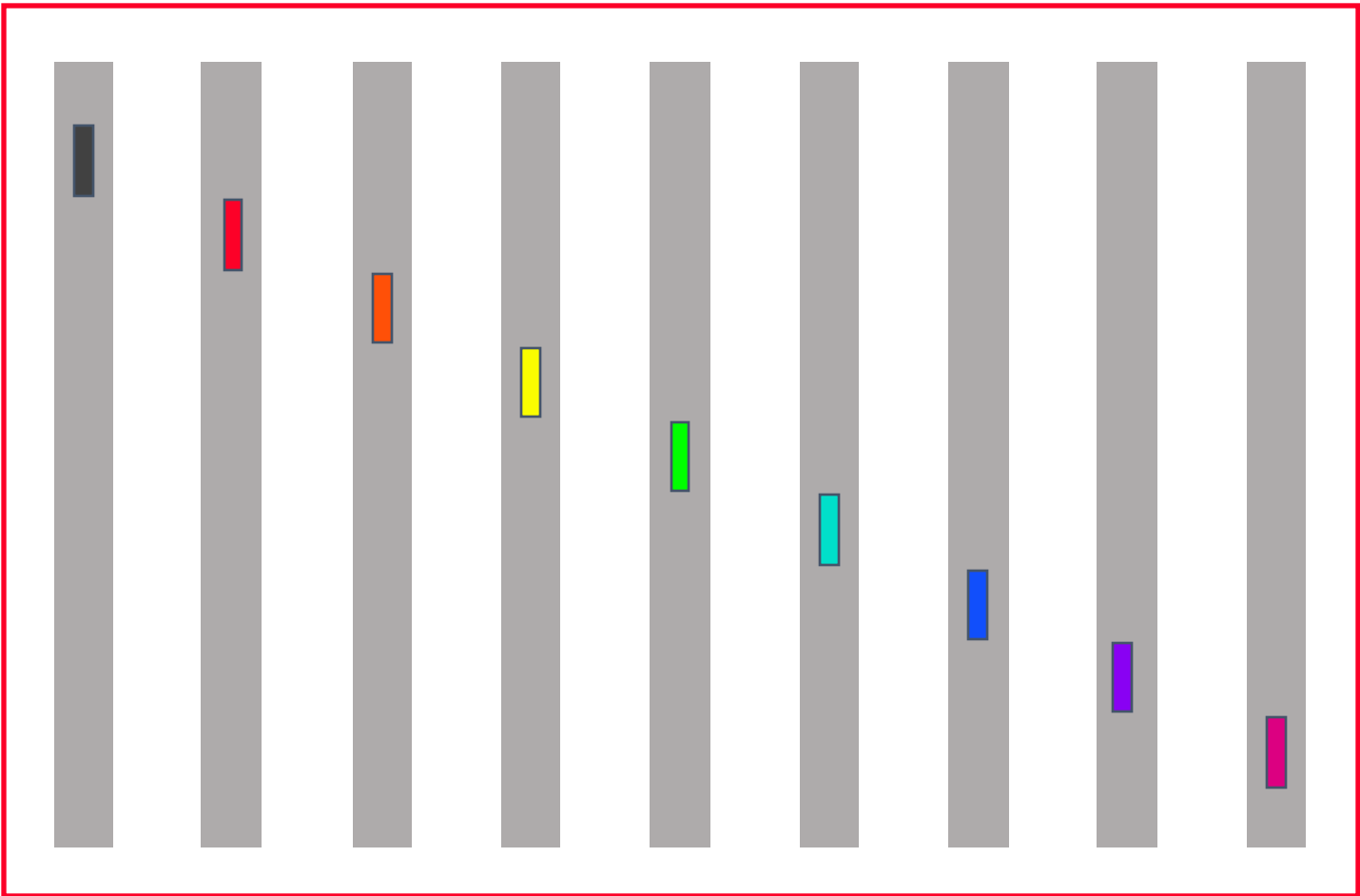
# Reduce(-to-one) (Large Message)



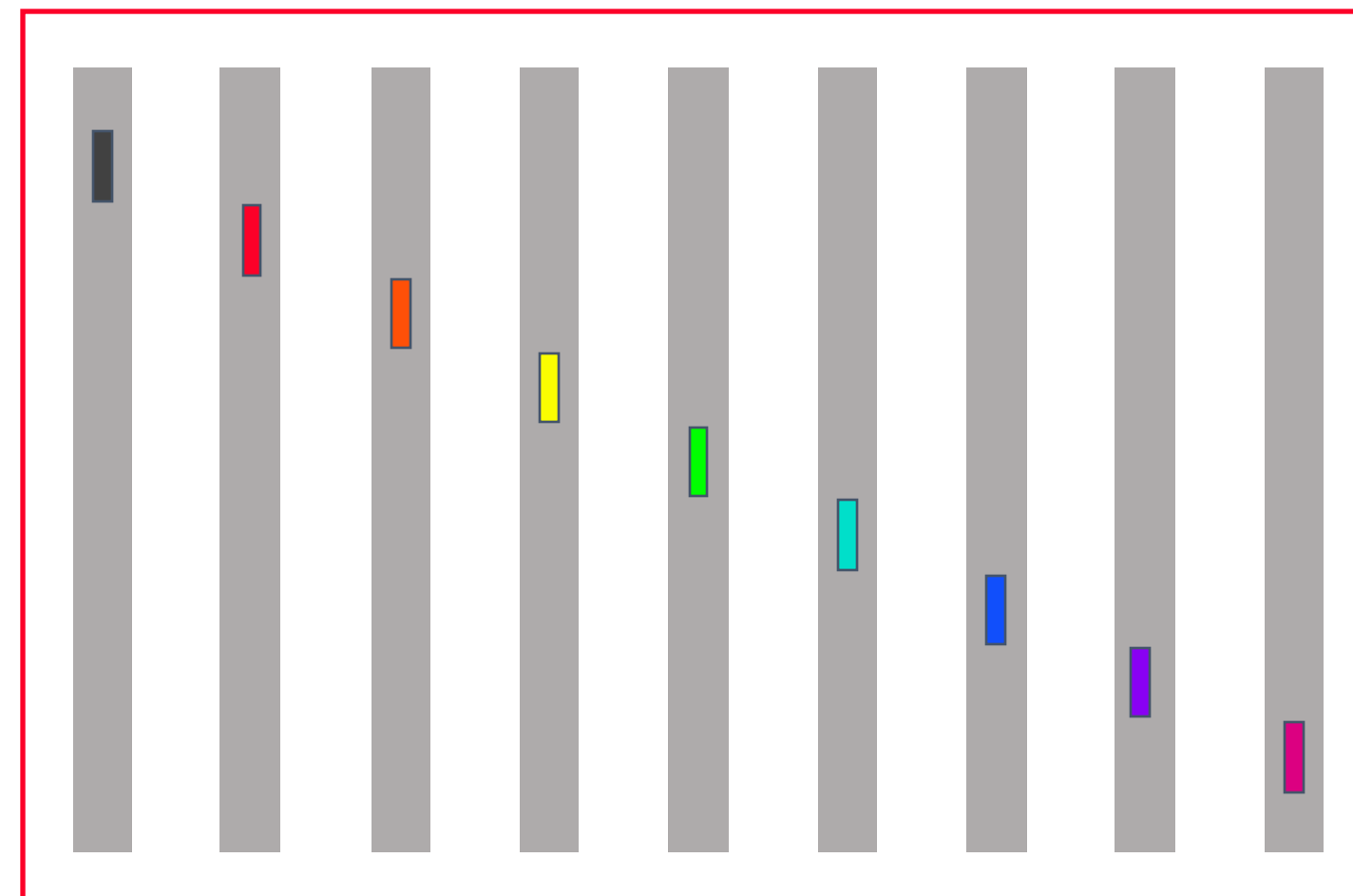
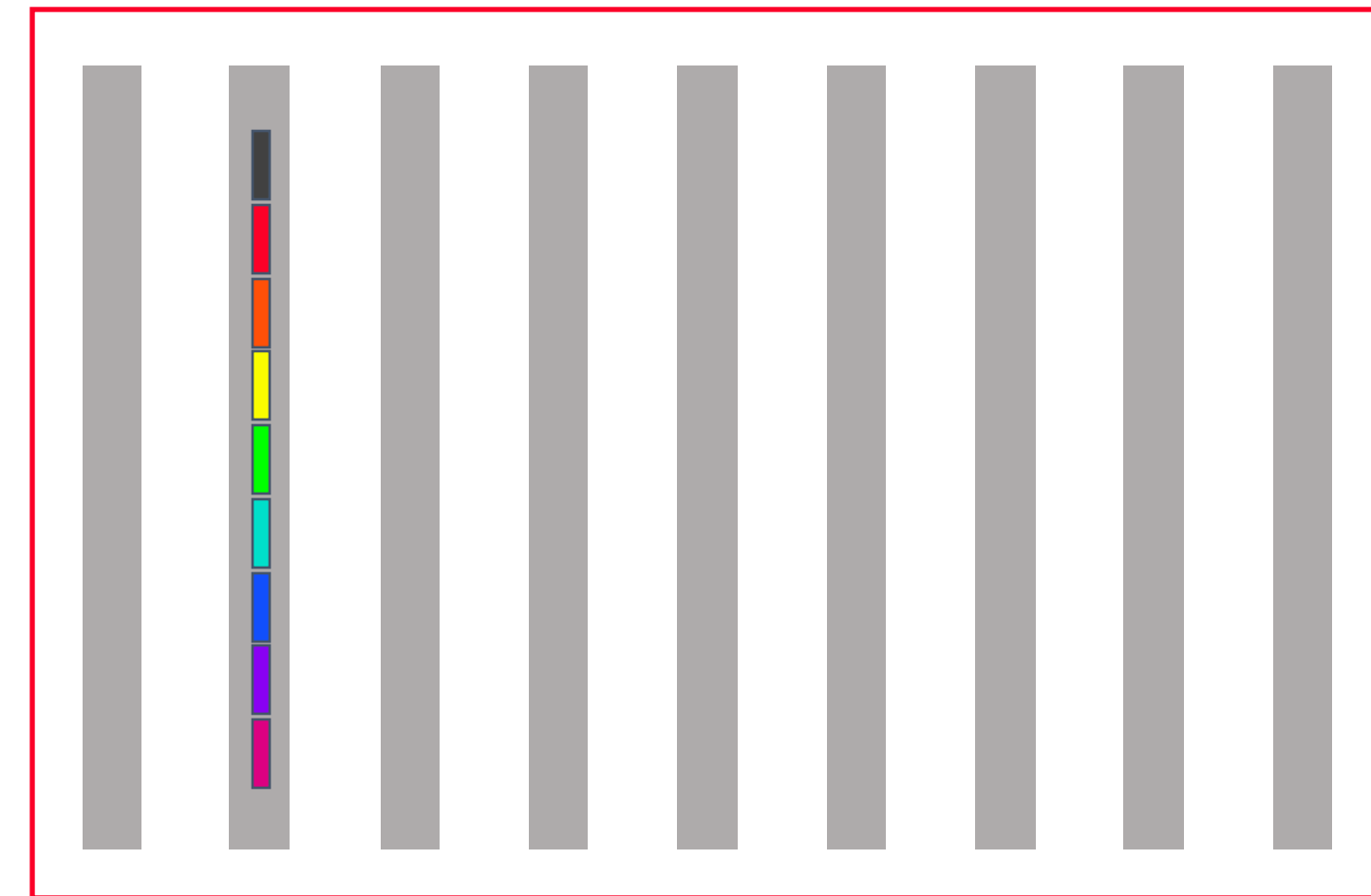
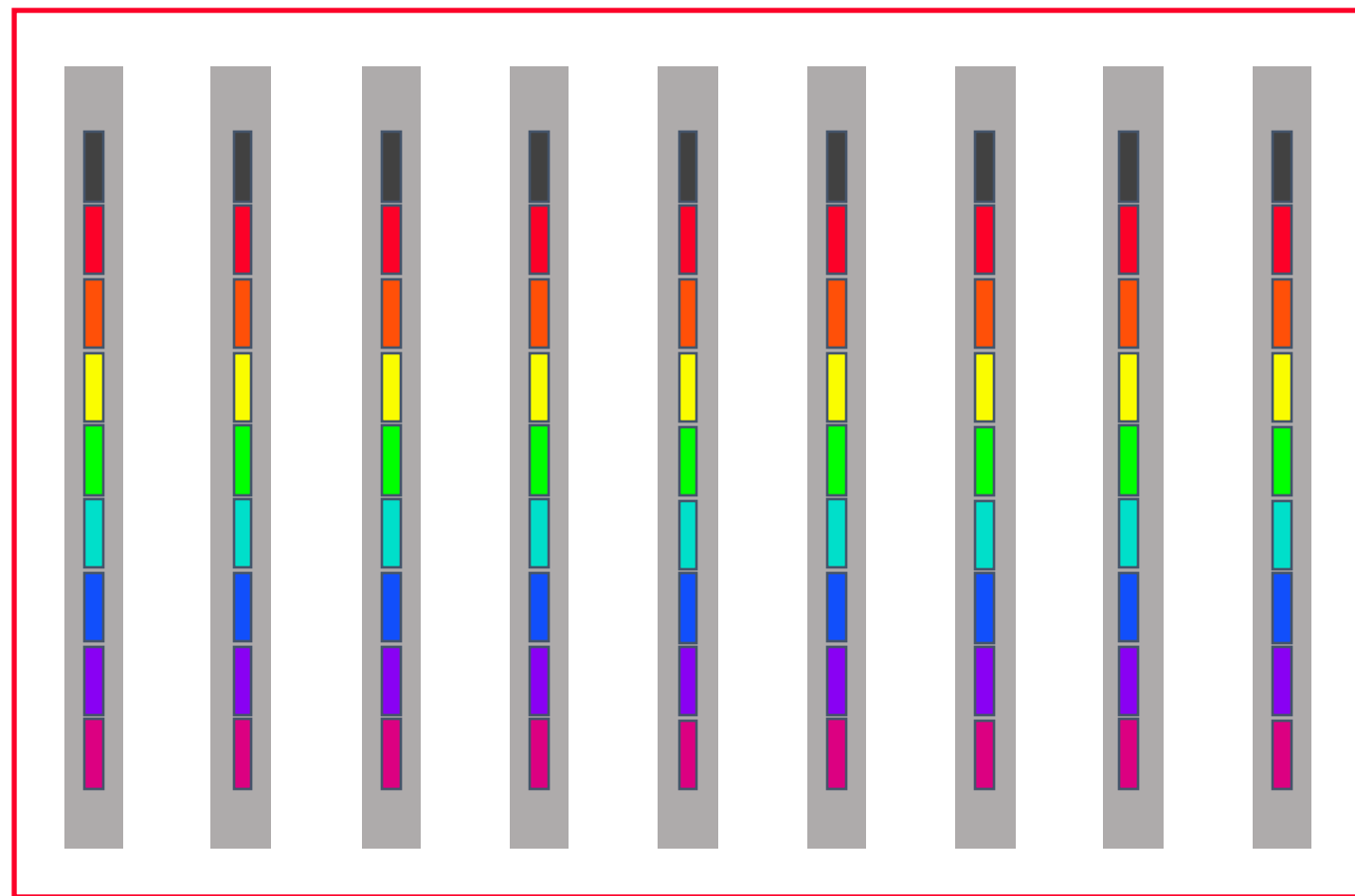
# Reduce (long vector)



Reduce-scatter



# Combine-to-one (long vector)



Gather

# Cost of Reduce-scatter/Gather Reduce(-to-one)

- Assumption: power of two number of nodes

Reduce-scatter  $(p-1)\alpha + \frac{p-1}{p}n\beta + \frac{p-1}{p}n\gamma$

gather  $\log(p)\alpha + \frac{p-1}{p}n\beta$

---

$(\log(p) + p - 1)\alpha + 2\frac{p-1}{p}n\beta + \frac{p-1}{p}n\gamma$

# Cost of Reduce-scatter/Gather Reduce(-to-one)

- Assumption: power of two number of nodes

Reduce-scatter  $(p-1)\alpha + \frac{p-1}{p}n\beta + \frac{p-1}{p}n\gamma$

gather  $\log(p)\alpha + \frac{p-1}{p}n\beta$

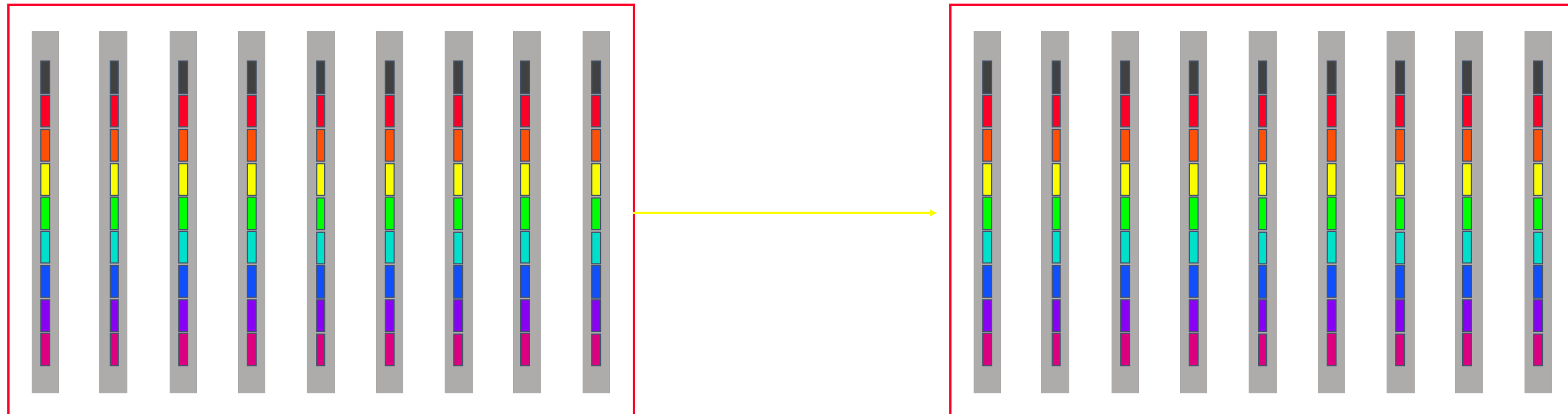
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$$(\log(p) + p - 1)\alpha + 2\frac{p-1}{p}n\beta + \frac{p-1}{p}n\gamma$$

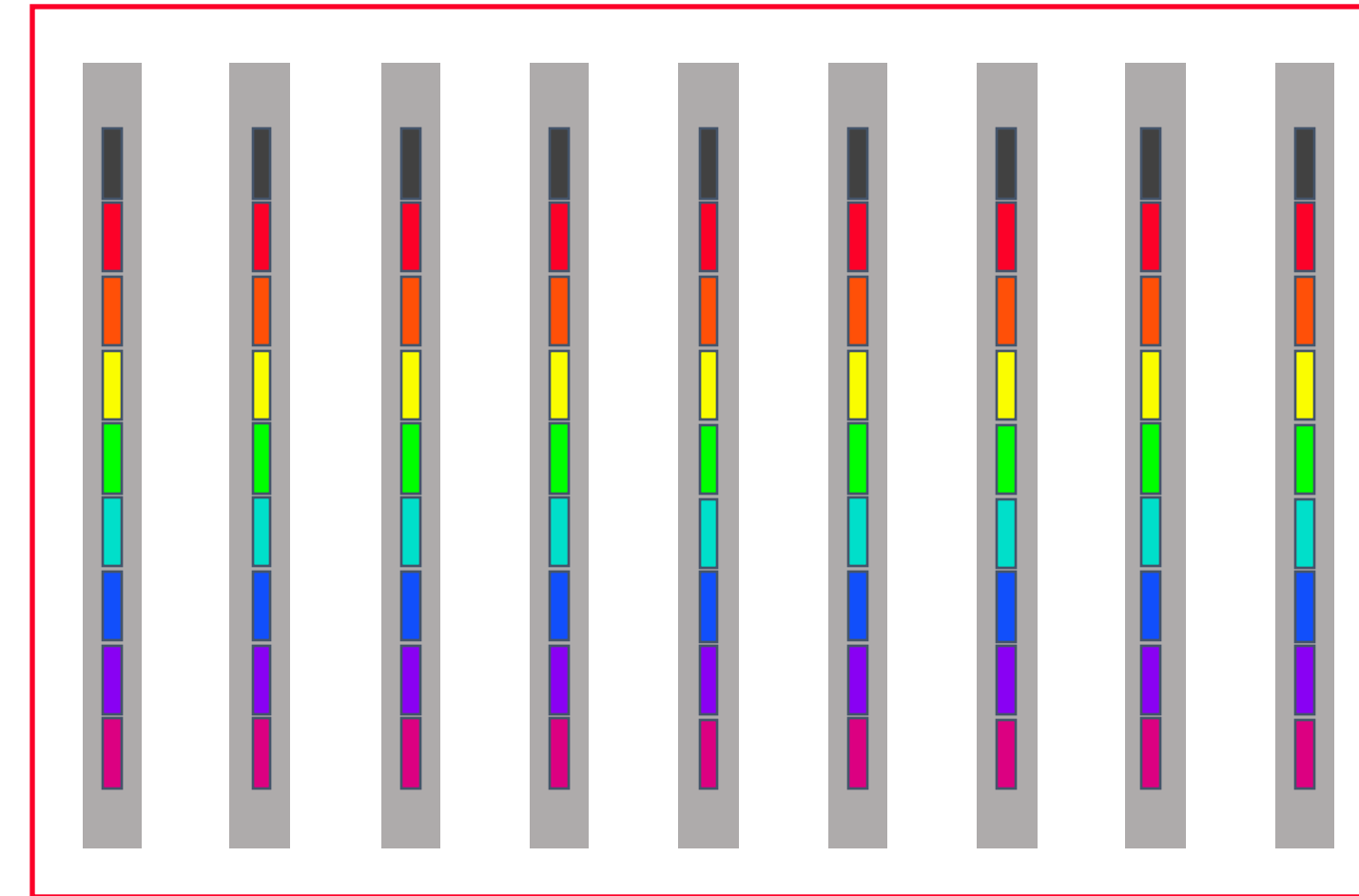
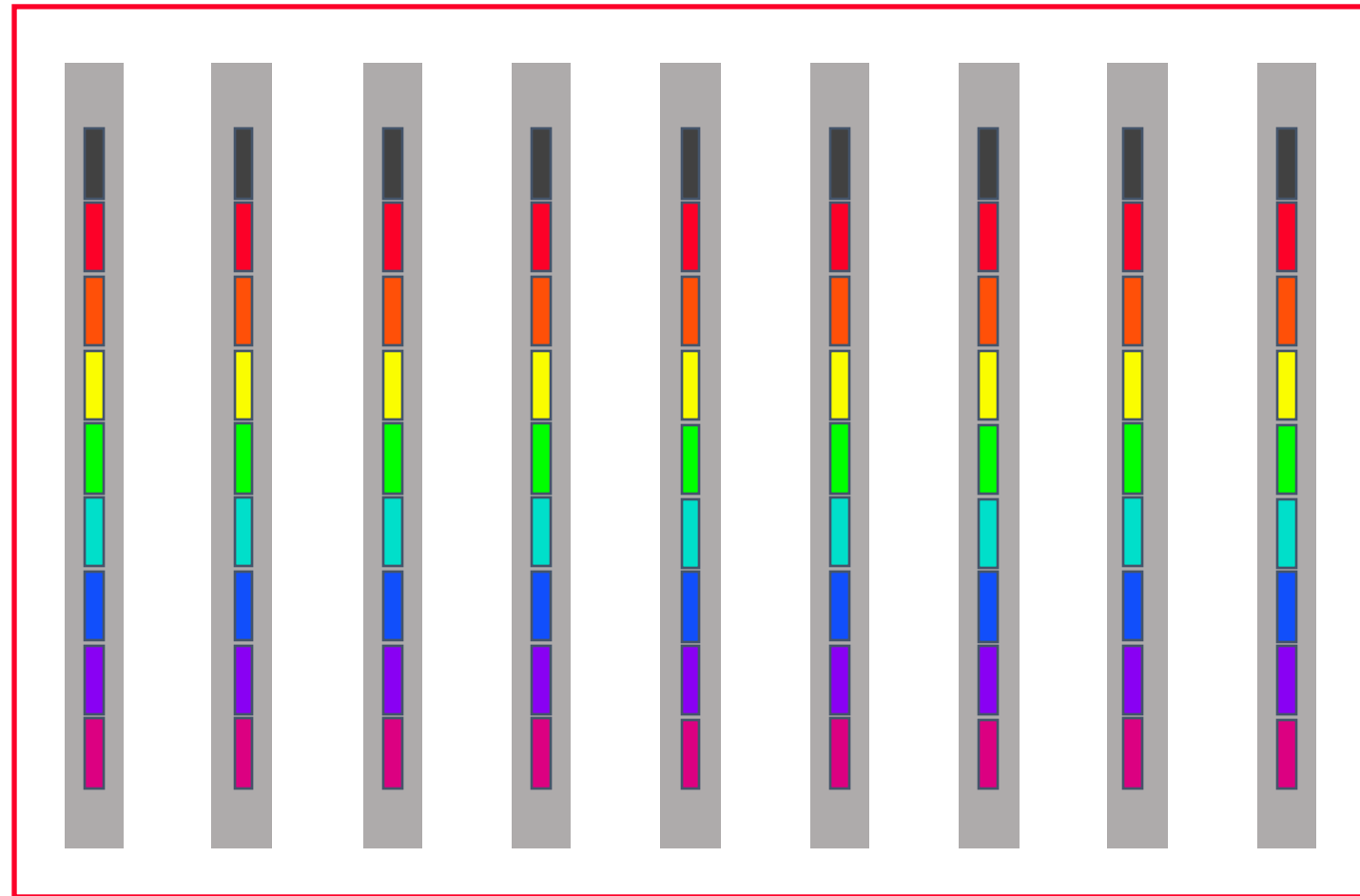
Vs. MST reduce:  $\lceil \log(p) \rceil (\alpha + n\beta + n\gamma)$



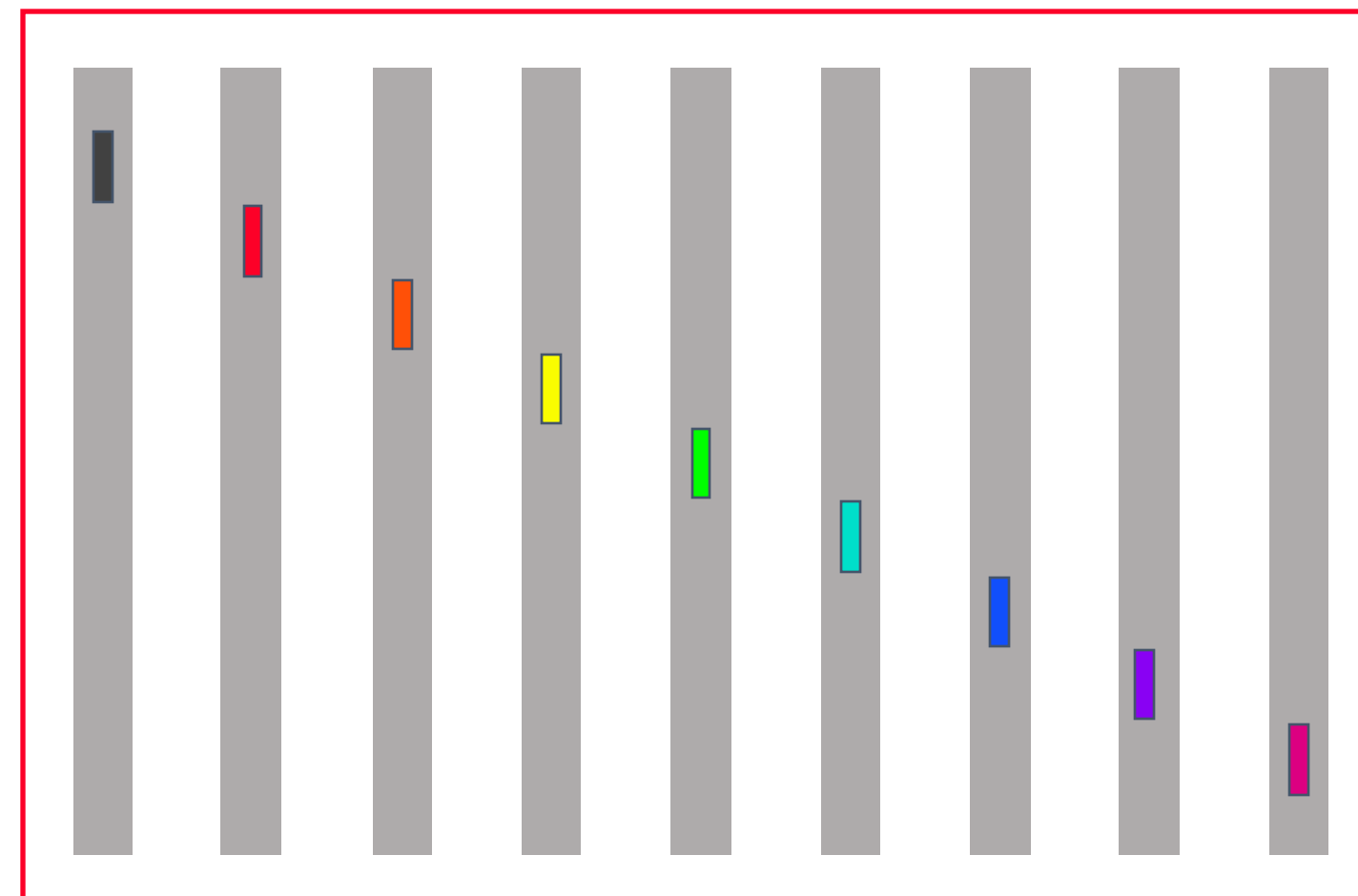
# Allreduce (Large Message)



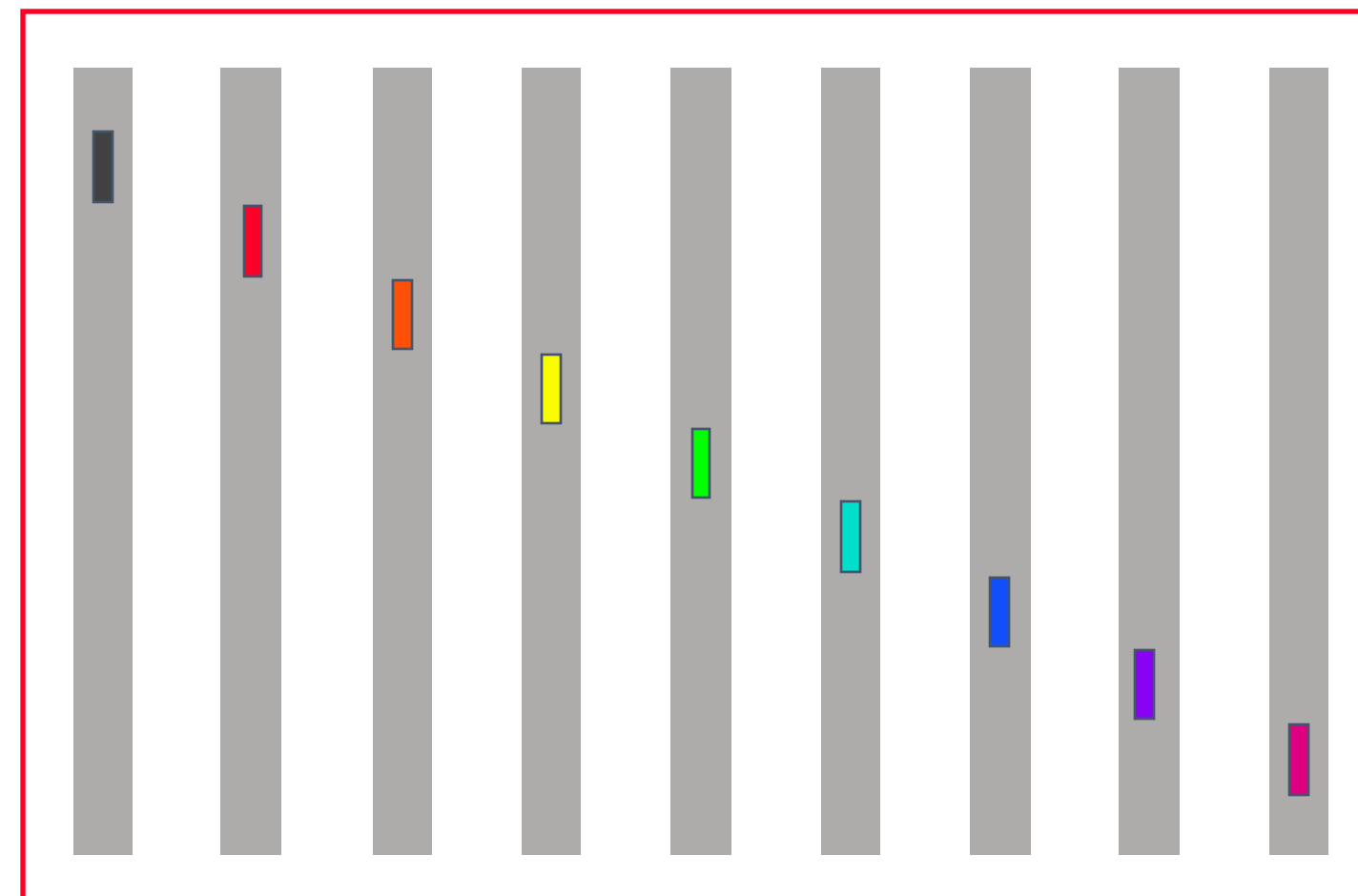
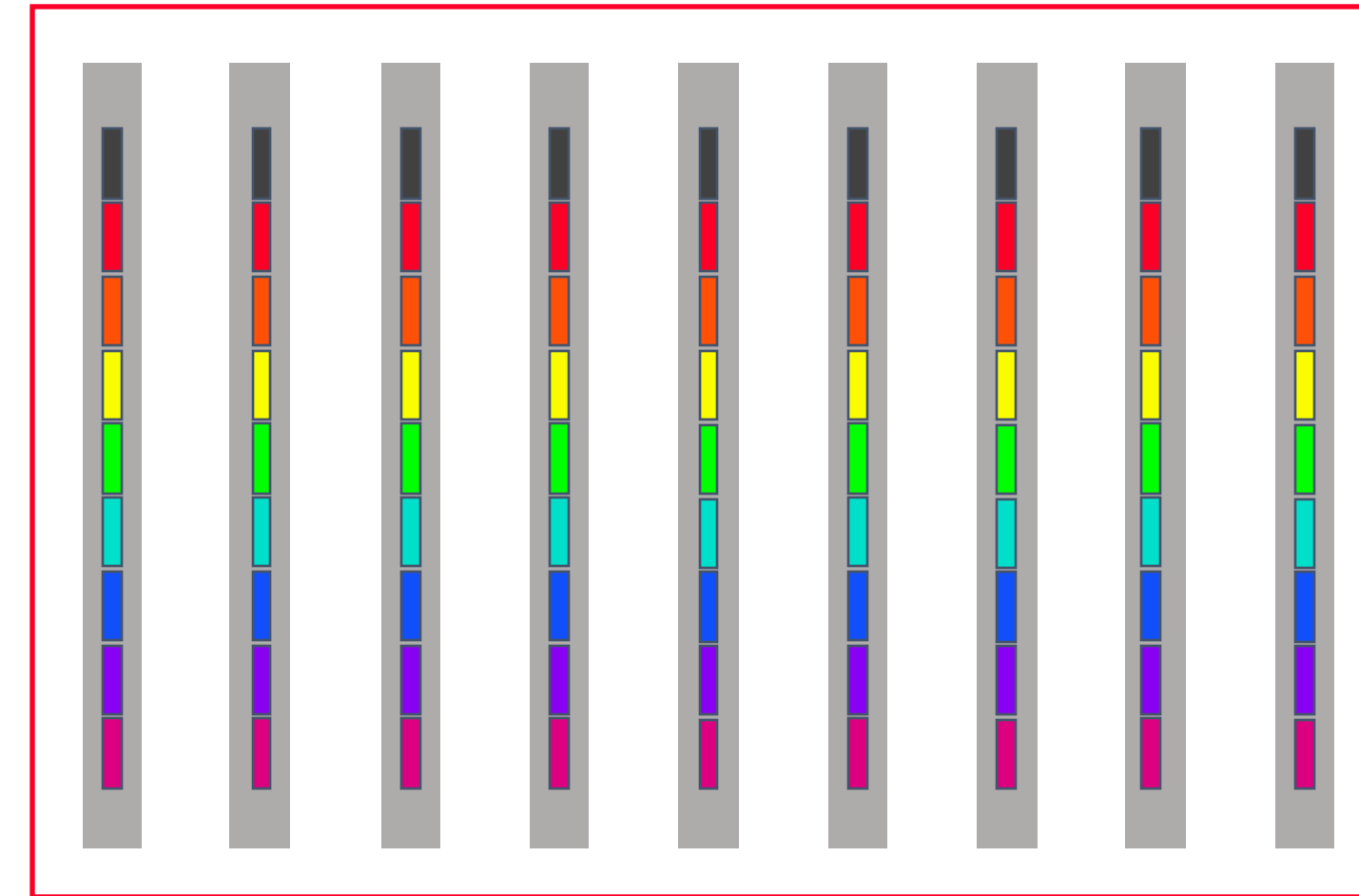
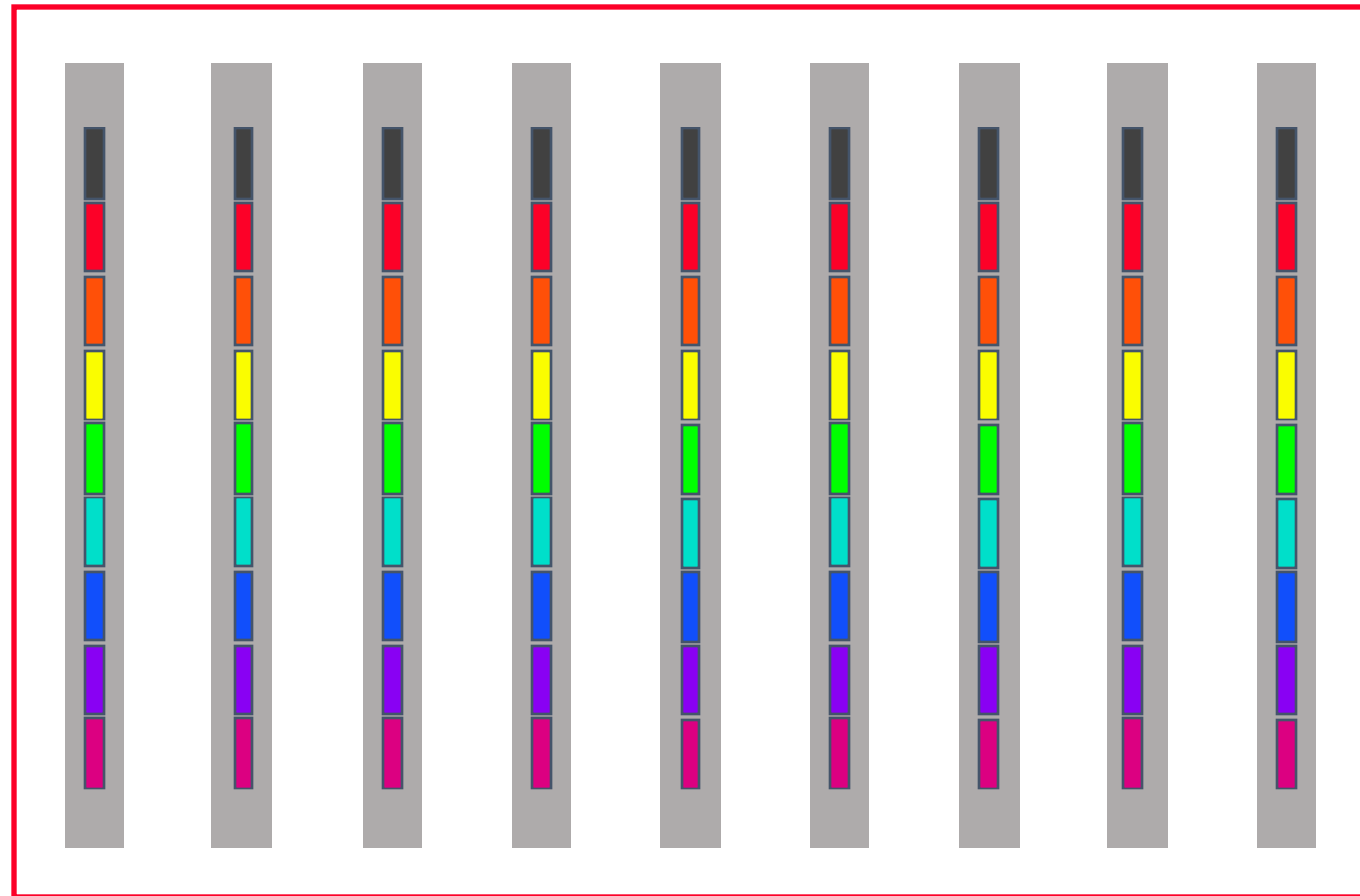
# Allreduce (Large Message)



Reduce-scatter



# Allreduce (long vector)



Allgather

# Cost of Reduce-scatter/Allgather Allreduce

- Assumption: power of two number of nodes

$$\text{Reduce-scatter} \quad (p-1)\alpha + \frac{p-1}{p}n\beta + \frac{p-1}{p}n\gamma$$

$$\text{Allgather} \quad \frac{(p-1)\alpha + \frac{p-1}{p}n\beta}{2(p-1)\alpha + 2\frac{p-1}{p}n\beta + \frac{p-1}{p}n\gamma}$$

## Cost of Reduce-scatter/Allgather Allreduce

- Assumption: power of two number of nodes

Reduce-scatter  $(p-1)\alpha + \frac{p-1}{p}n\beta + \frac{p-1}{p}n\gamma$

Allgather  $(p-1)\alpha + \frac{p-1}{p}n\beta$

---

$2(p-1)\alpha + 2\frac{p-1}{p}n\beta + \frac{p-1}{p}n\gamma$

Vs. Reduce-broadcast  
allreduce  $2\log(p)\alpha + 2\log(p)n\beta + \log(p)n\gamma$

Recap

Reduce-scatter

$$(p-1)\alpha + \frac{p-1}{p}n(\beta + \gamma)$$

Scatter

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Gather

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Allgather

$$(p-1)\alpha + \frac{p-1}{p}n\beta$$

Reduce(-to-one)

Allreduce

Broadcast

# Recap

Reduce-scatter

$$(p-1)\alpha + \frac{p-1}{p}n(\beta + \gamma)$$

Scatter

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Gather

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Allgather

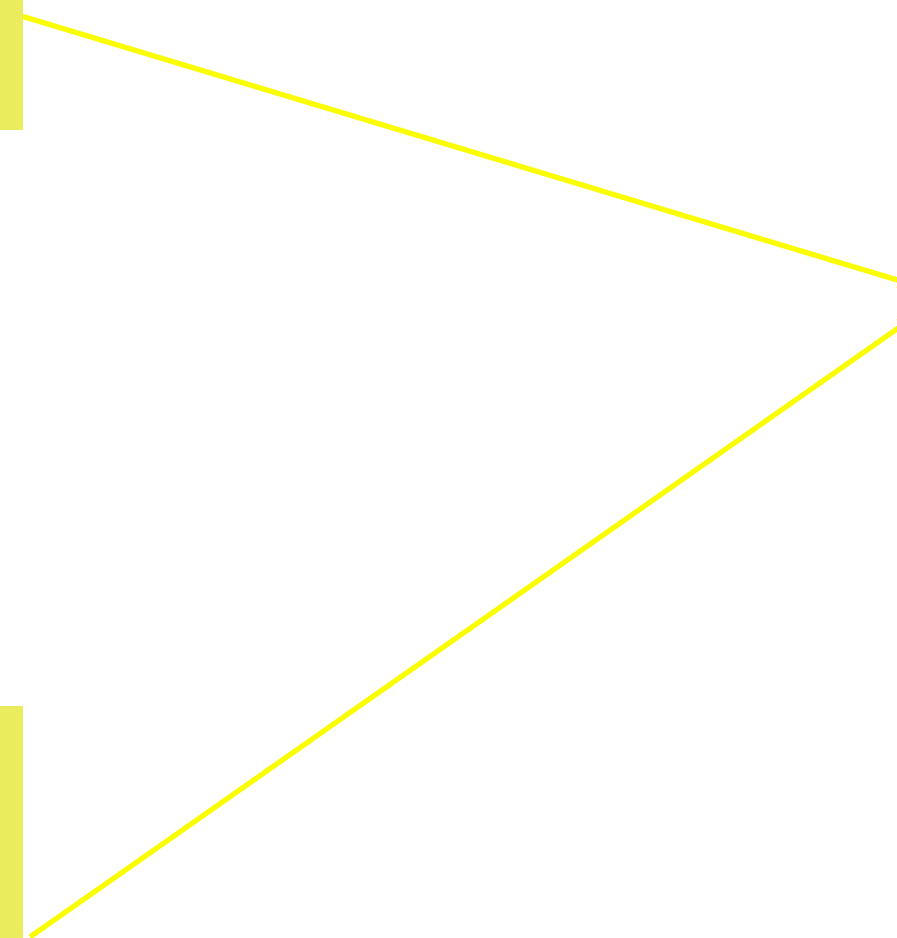
$$(p-1)\alpha + \frac{p-1}{p}n\beta$$

Reduce(-to-one)

$$(p-1 + \log(p))\alpha + \frac{p-1}{p}n(2\beta + \gamma)$$

Allreduce

Broadcast



# Recap

## Reduce-scatter

$$(p-1)\alpha + \frac{p-1}{p}n(\beta + \gamma)$$

### Scatter

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

### Gather

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

### Allgather

$$(p-1)\alpha + \frac{p-1}{p}n\beta$$

## Reduce(-to-one)

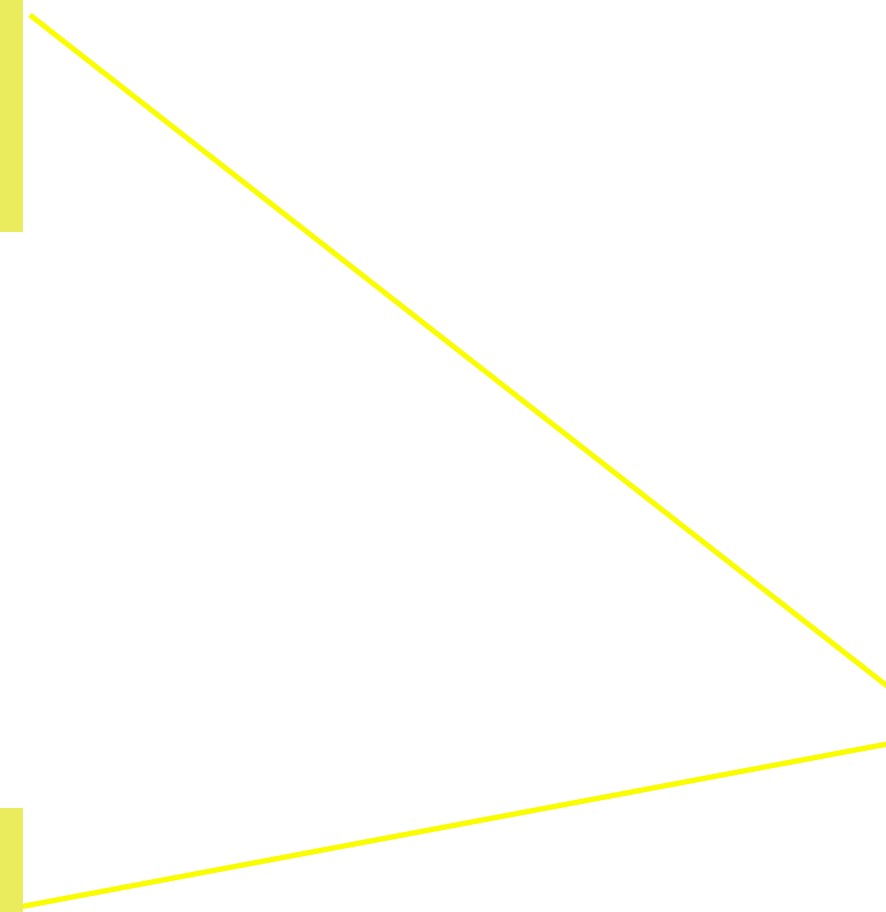
$$(p-1 + \log(p))\alpha + \frac{p-1}{p}n(2\beta + \gamma)$$

## Allreduce

$$2(p-1)\alpha + \frac{p-1}{p}n(2\beta + \gamma)$$

## Broadcast

$$(\log(p) + p-1)\alpha + 2\frac{p-1}{p}n\beta$$





# Recap

Reduce-scatter

$$(p-1)\alpha + \frac{p-1}{p}n(\beta + \gamma)$$

Scatter

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Gather

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Allgather

$$(p-1)\alpha + \frac{p-1}{p}n\beta$$

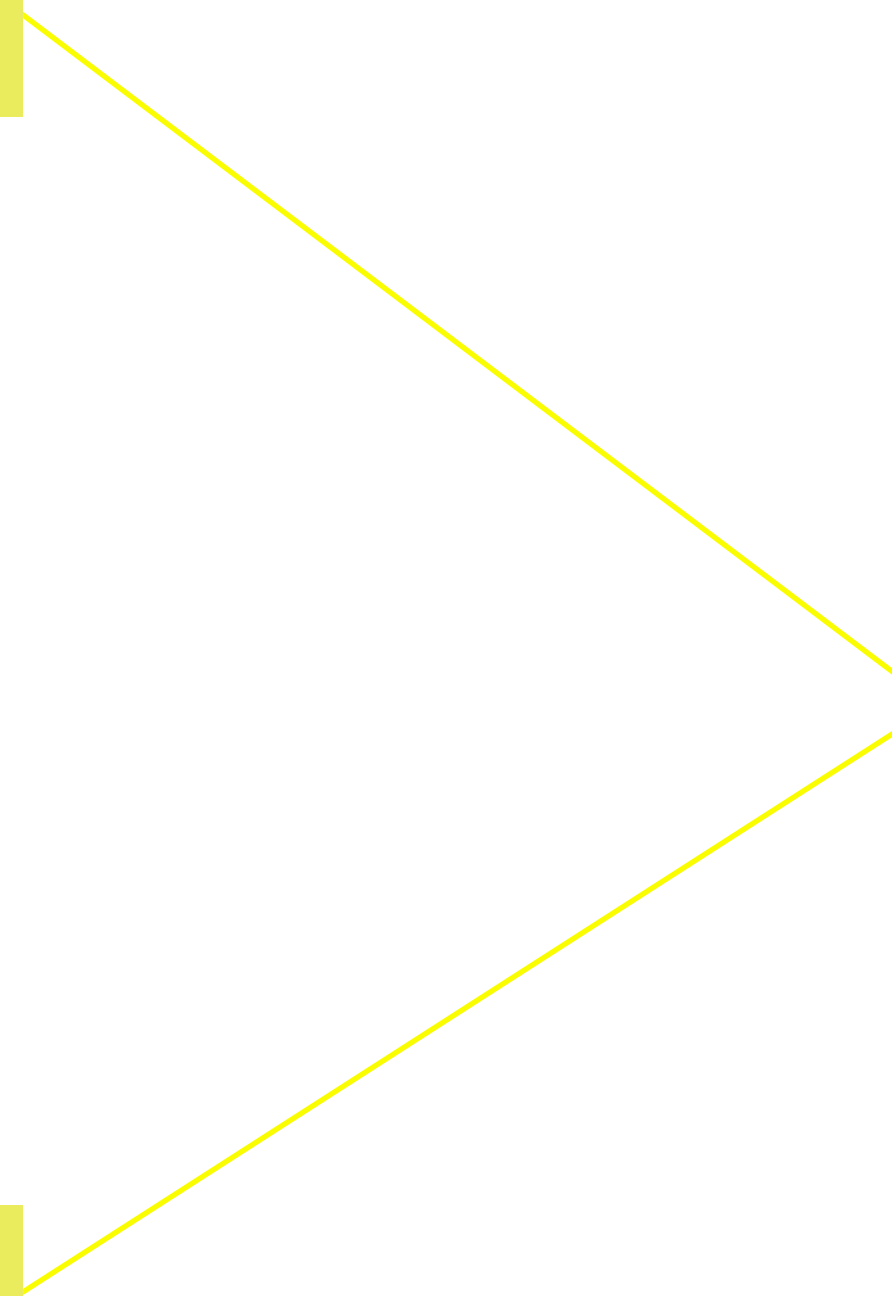
Reduce(-to-one)

$$(p-1+\log(p))\alpha + \frac{p-1}{p}n(2\beta + \gamma)$$

Allreduce

$$2(p-1)\alpha + \frac{p-1}{p}n(2\beta + \gamma)$$

Broadcast



# Recap

Reduce-scatter

$$(p-1)\alpha + \frac{p-1}{p}n(\beta + \gamma)$$

Scatter

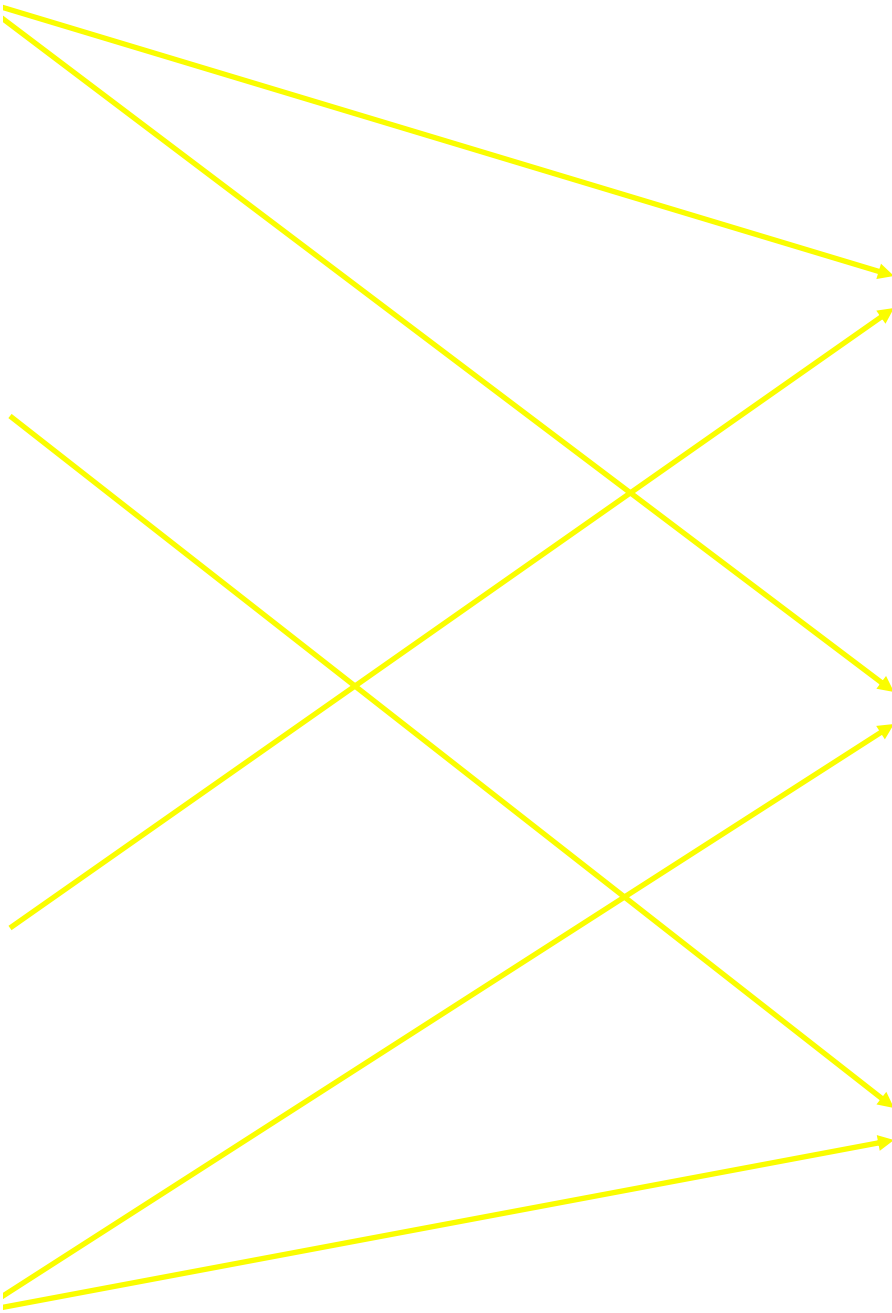
$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Gather

$$\log(p)\alpha + \frac{p-1}{p}n\beta$$

Allgather

$$(p-1)\alpha + \frac{p-1}{p}n\beta$$



Reduce(-to-one)

$$(p-1+\log(p))\alpha + \frac{p-1}{p}n(2\beta + \gamma)$$

Allreduce

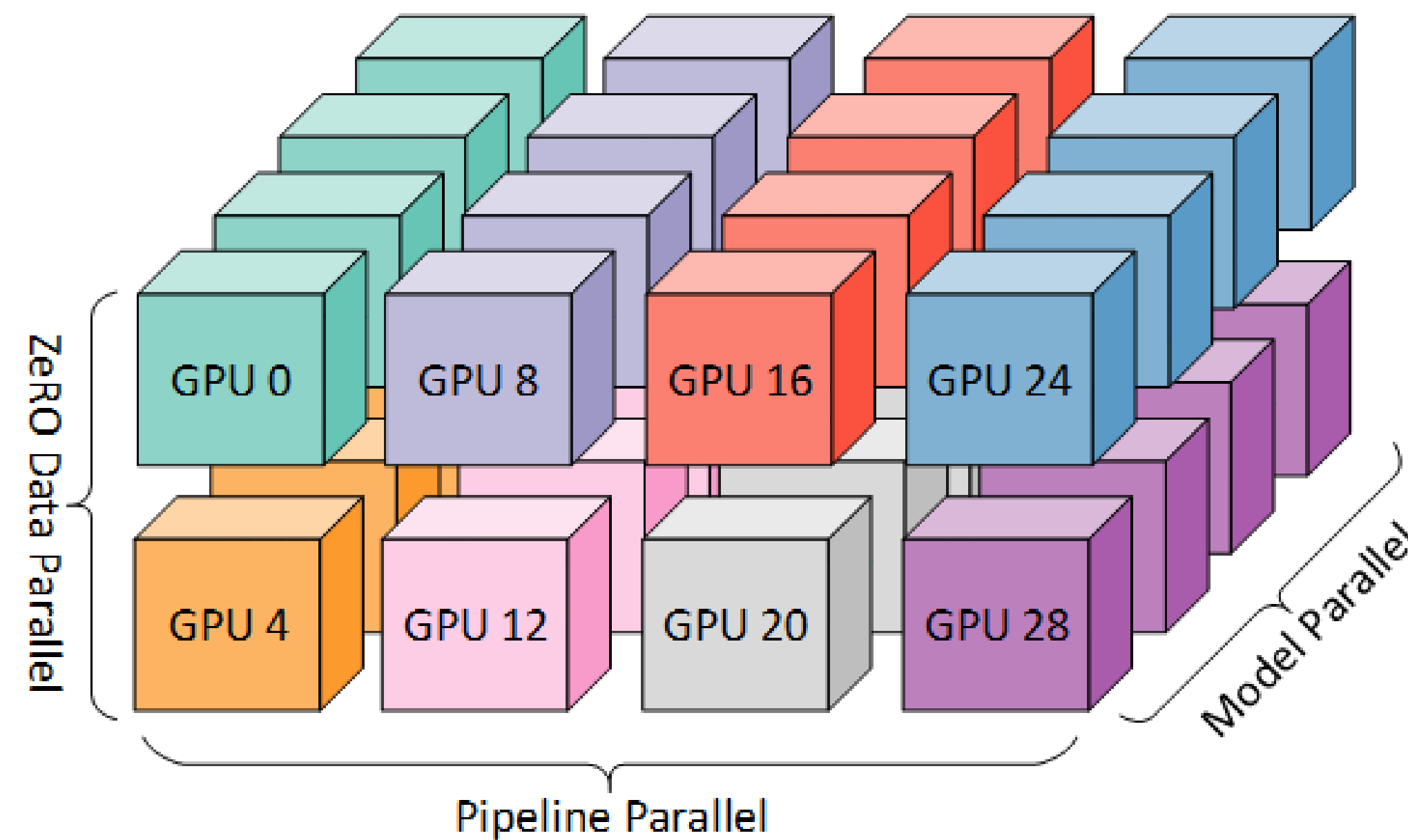
$$2(p-1)\alpha + \frac{p-1}{p}n(2\beta + \gamma)$$

Broadcast

$$(\log(p) + p-1)\alpha + 2\frac{p-1}{p}n\beta$$

# A More Complicate Case

- Real Cluster to train ChatGPT:
  - If using GPU: 2D Mesh
  - If using TPU: 3D Mesh, see figure below



# Summary and Question

- MST -> when  $\alpha$  dominates
- Ring -> when  $n \cdot \beta$  dominates
- 2D can be composed using 1D, 3D can be composed using 2D,  
...
- Latency / Bandwidth trade-offs

# Recap

- Q1: Which collective primitive maps to the distributed SGD gradient synchronization step?
- Q2: How many messages do we need to transfer over the network for a single iteration of GPT-3 SGD update assuming 8-gpu parallelism?
- Q3: For Q2, assuming 1D mesh, should we use MST or Ring?

# Collective Pros

- A set of structured / well-defined communication primitives
- Extremely well-optimized
- Beautiful math, easy to analyze, and easy to understand its performance

# Collective Cons

- Lack of Fault Tolerance
  - What if one node (in the ring) is dead?
- Requires Homogeneity
  - What if one node computes slower than all other nodes?
  - What if one link has lower bandwidth than the other node?

## Real Cluster:

- Need Fault tolerance
- Heterogeneous hardware setup

# Where we are

Motivations, Economics, Ecosystems,  
Trends



Skip this  
Storage

## Networking

## Compute

Datacenter  
networking

Collective  
communication

(Distributed) File  
Systems / Database

Cloud storage

Distributed  
Computing

Big data  
processing



# But Some Basic Knowledge Check





- **What is a Database**
- Have you heard these terms?
  - HashTable
  - SSTable and LSM-Trees
  - BTree?
- Optional readings will cover this – highly recommend to read

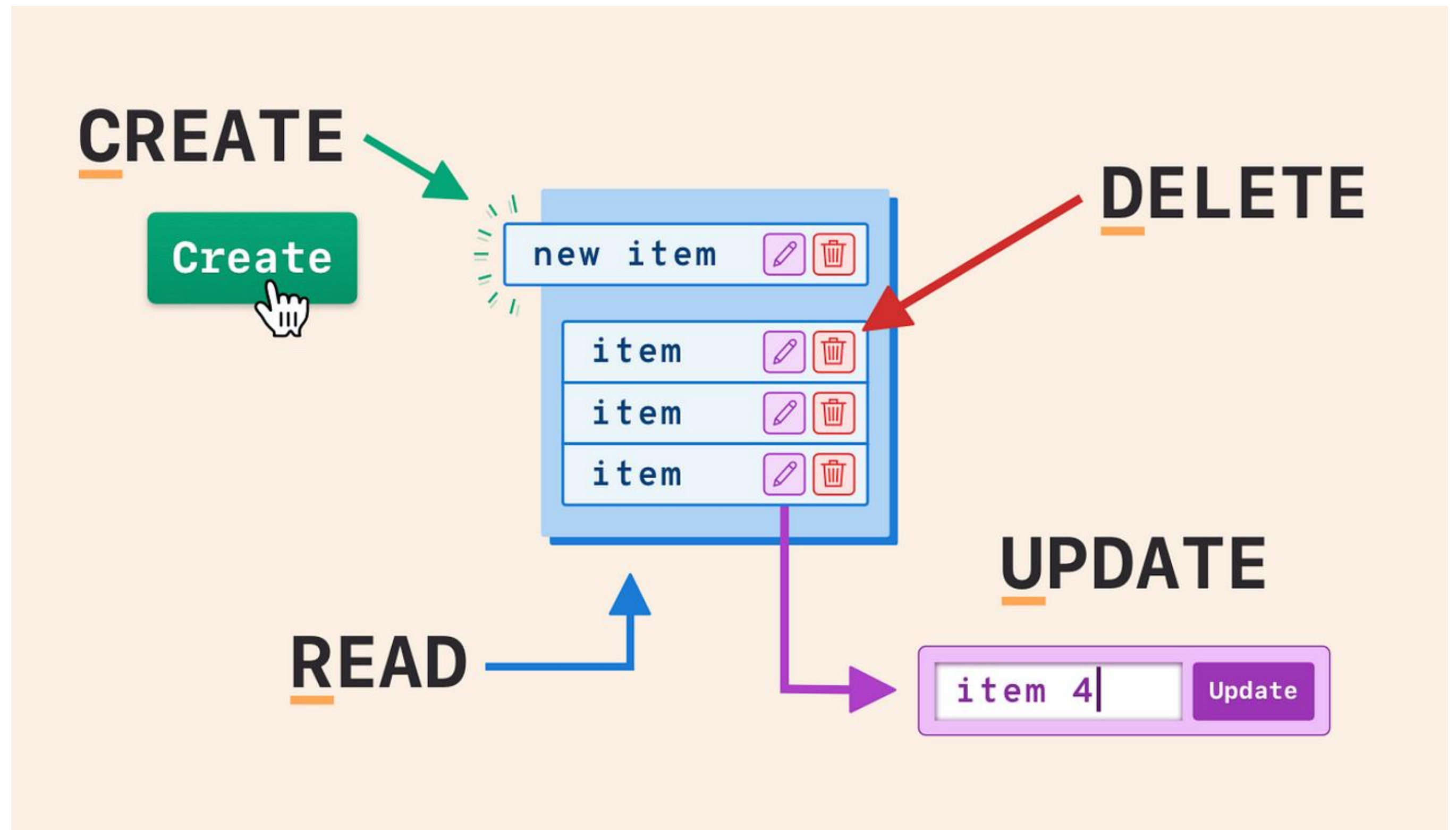
# But Some Important Concepts

- OLTP v.s. OLAP
- Data warehousing
- Schemas for Analytics
- Column-oriented storage
- Data cubes and materialized views

# CRUD

I'm a Database Developer,  
all what I do is

	<b><i>C</i></b> reate
	<b><i>R</i></b> ead
	<b><i>U</i></b> pdate
	<b><i>D</i></b> elele



# Database transactions

- Make sale
- Place an order
- Pay an employee's salary
- Comment a blog post
- Act in games
- Add/remove contract to an address book

Online transaction processing (OLTP)

# Walmart Beer and Diaper (1988)



- Unexpected correlation:
  - Sales of diapers and beer

Forbes 1988

# Data analytics

- What was the total revenue of each of our stores in Jan?
- How many more bananas than usual did we sell during our latest data?
- Which brand of baby food is most often purchased together with brand X diapers?

Online analytic processing (OLAP)

# OLTP v.s. OLAP

Property	Transaction processing systems (OLTP)	Analytic systems (OLAP)
Main read pattern	Small number of records per query, fetched by key	Aggregate over large number of records



# OLTP v.s. OLAP

Property	Transaction processing systems (OLTP)	Analytic systems (OLAP)
Main read pattern	Small number of records per query, fetched by key	Aggregate over large number of records
Main write pattern	Random-access, low-latency writes from user input	Bulk import (ETL) or event stream
Primarily used by	End user/customer, via web application	Internal analyst, for decision support
What data represents	Latest state of data (current point in time)	History of events that happened over time
Dataset size	Gigabytes to terabytes	Terabytes to petabytes



# Today's topic

- OLTP v.s. OLAP
- Data warehousing
- Schemas for Analytics
- Column-oriented storage



# Transaction systems are complex.

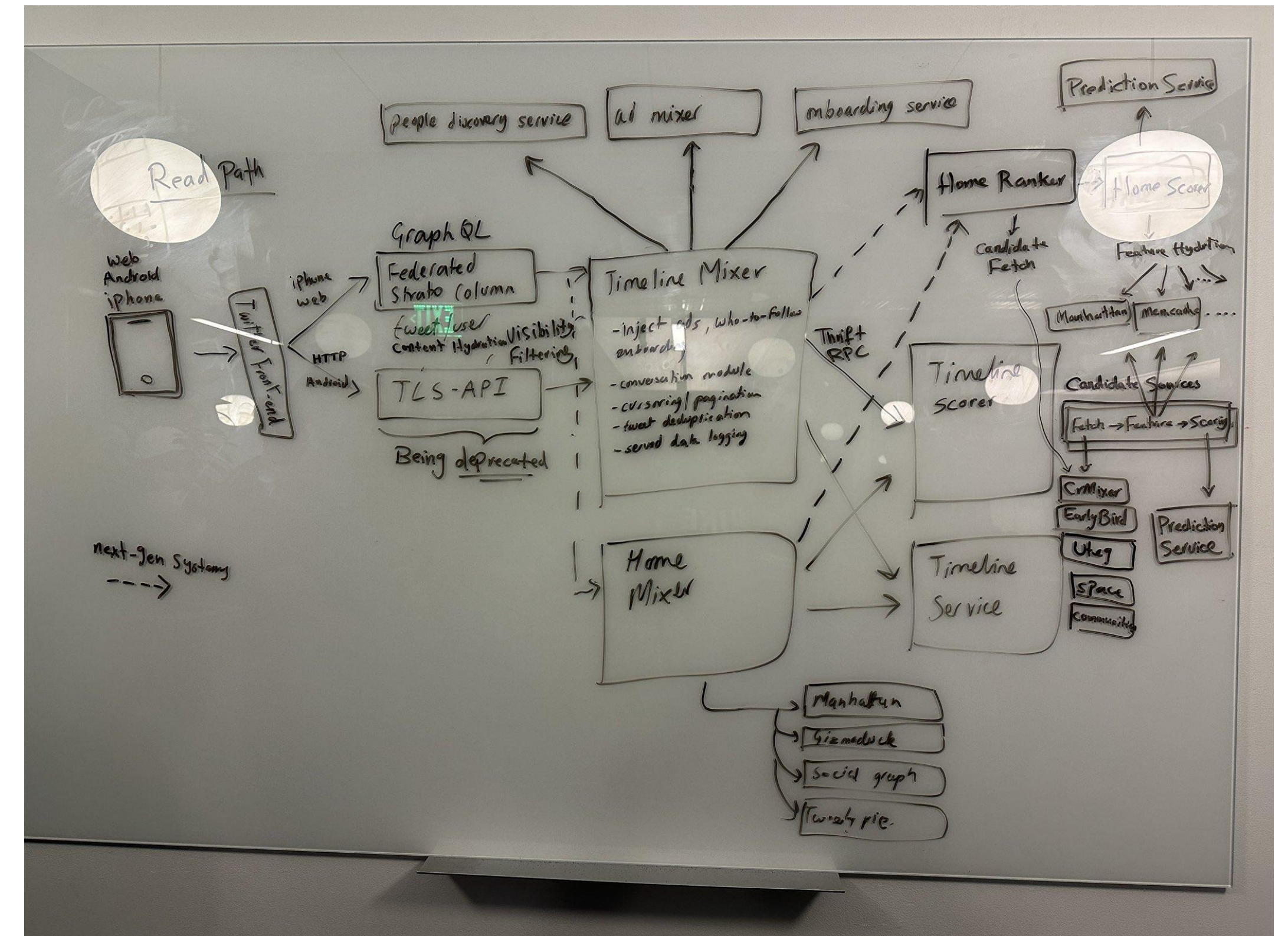


Elon Musk  
@elonmusk

Just leaving Twitter HQ code review



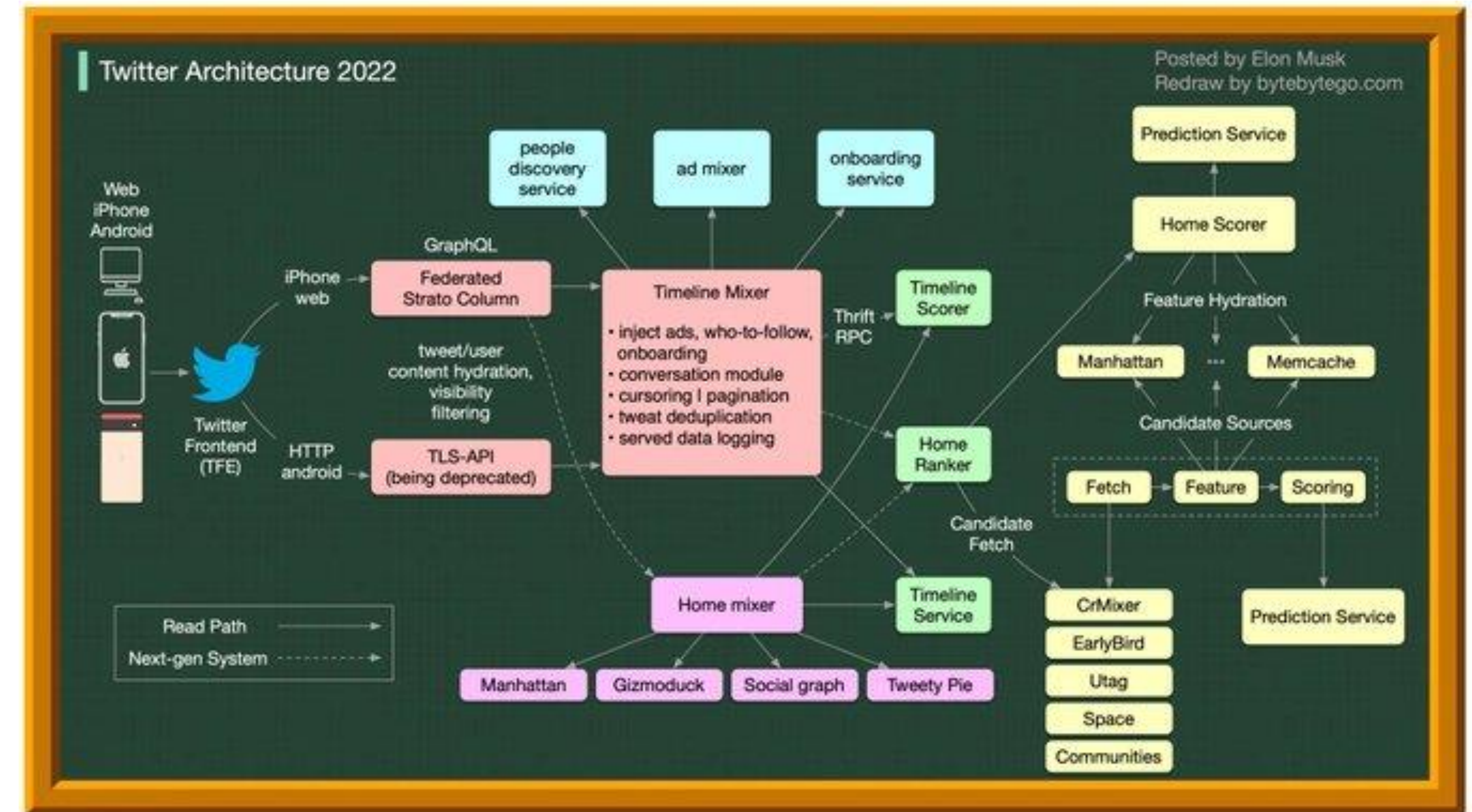
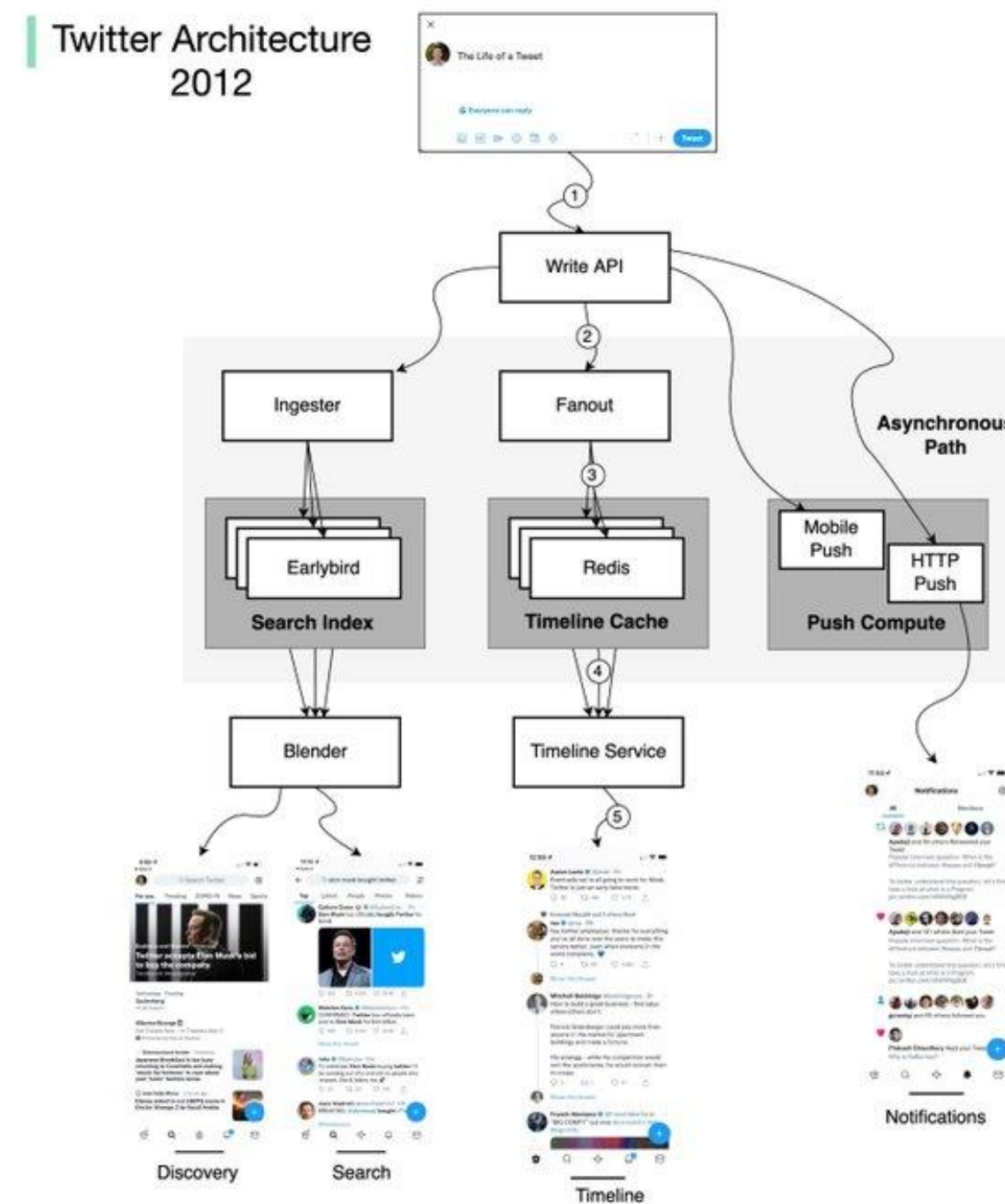
1:28 AM · Nov 19, 2022 · Twitter for iPhone



Elon Musk's Twitter System Design Diagram Explained  
[https://www.youtube.com/watch?v=\\_Y5aGCOkymQ](https://www.youtube.com/watch?v=_Y5aGCOkymQ)



# Transaction systems need to be highly available.



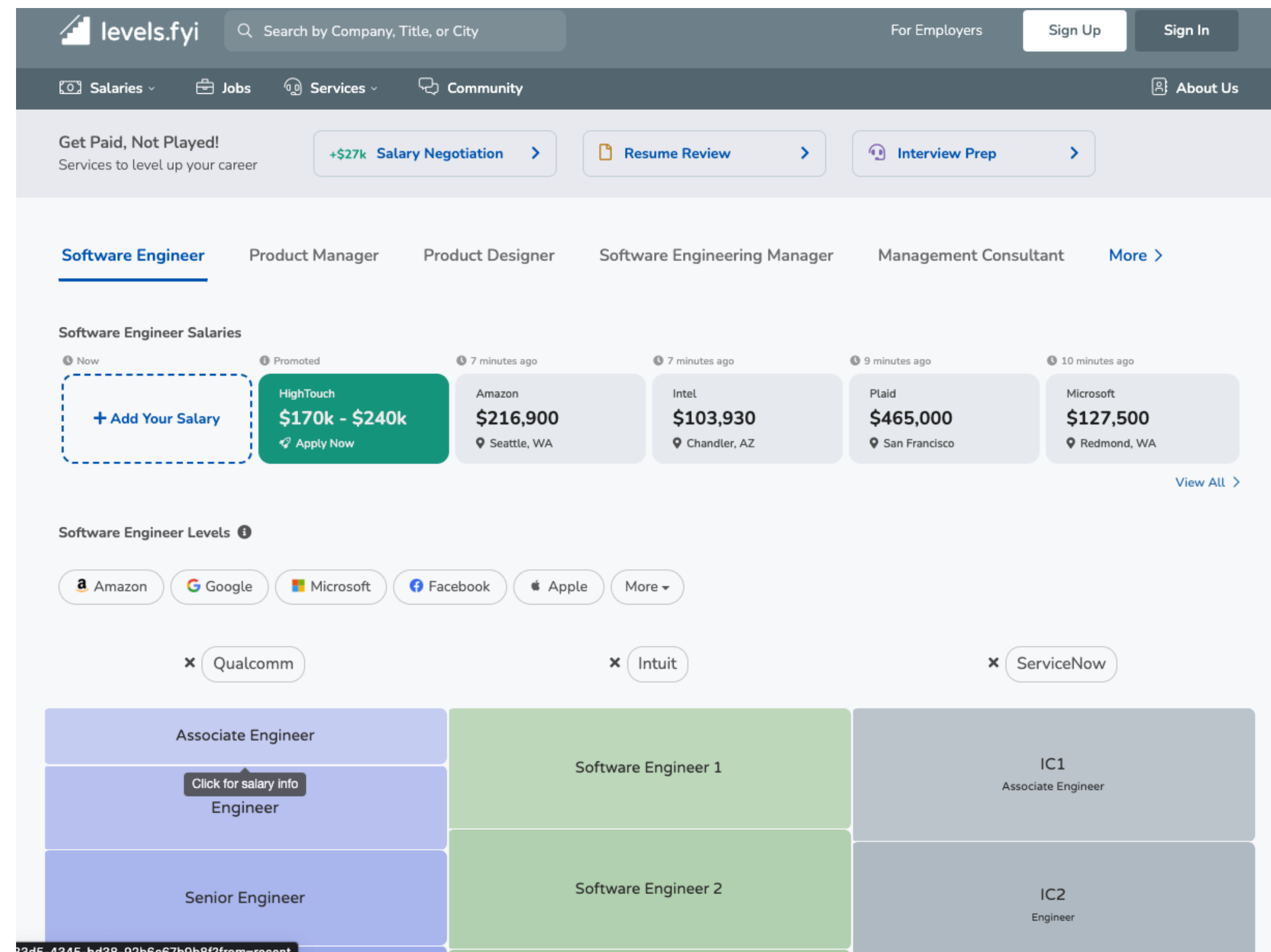
- Low latency.
- Highly available.
- Ad hoc analytic queries are expensive.

<https://twitter.com/alexxybyte/status/1594008281340530688>

# Data warehouse

- A separate database that analysts can query to their hearts' content, without affecting OLTP operations.
- Maintain a read-only copy for analytic purposes.
- Only exist in almost all large enterprises.

# Small companies?



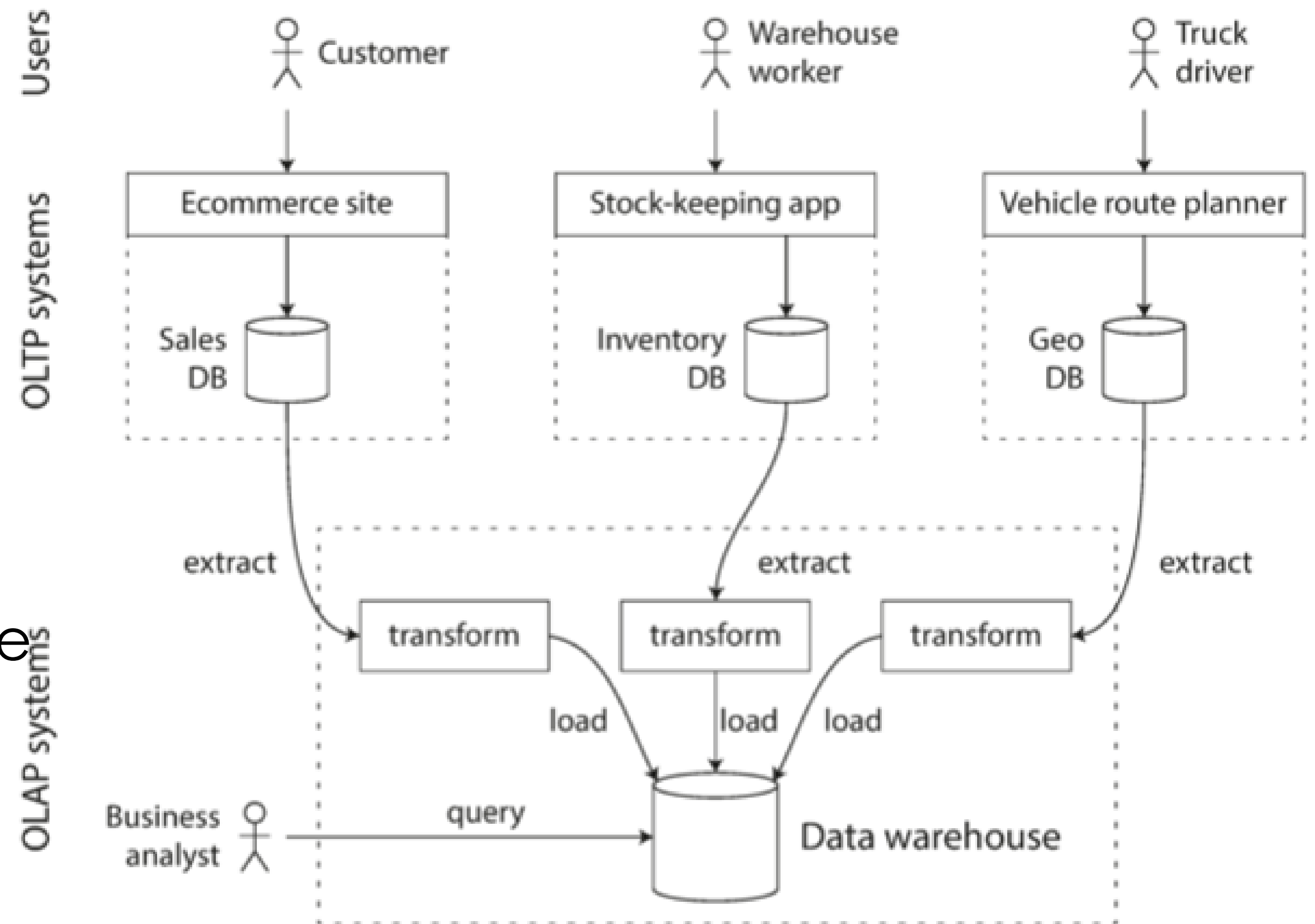
**How Levels.fyi scaled to millions of users with Google Sheets as a backend**

Our philosophy to scaling is simple, avoid premature optimization



# Extract-Transform-Load (ETL)

- Extract
  - Periodica data dump
  - Continuous streaming
- Transform
  - Analysis-friendly schema
  - Data cleaning
- Load into a data warehouse



# Why data warehouse?

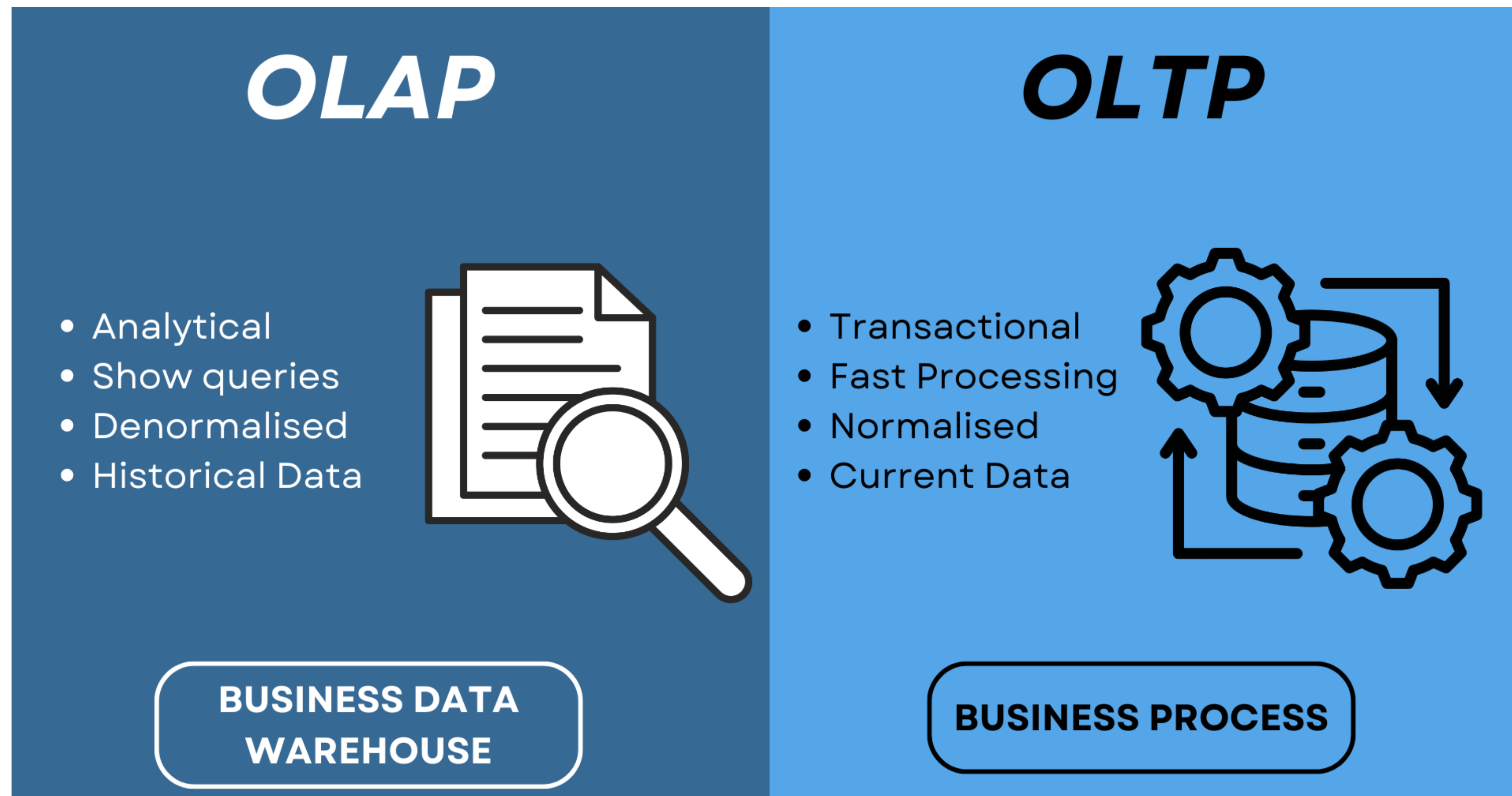
- Separation of concerns
  - Performance (reliability, latency)
  - Expertise requirement, management
- The classic indexes (e.g., SSTable, B-tree) are good for reading and writing a single record.
  - But are not good at answering analytic queries.

# How do you interact with OLAP & OLTP

- SQL query interface
  - *Select \* from*
  - “A database system can be considered mature when it has an SQL query interface”.
  - Both OLAP and OLTP
- OLAP:
  - More and more codeless user interfaces.
  - Text2SQL
  - Note: This is a big market of innovations



# Summary



More Stories?



400M



100B



800M / ~30 persons



83B

# Where We Are

Motivations, Economics, Ecosystems,  
Trends



Networking

Storage

**Part3: Compute**

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processing

# Where We Are

Machine Learning Systems

Big Data

Cloud

Foundations of Data Systems



2010 - Now

2000 - 2016

1980 - 2000

# Distributed Computing and Big Data

- Parallelism Basics
- Data Replication and partitioning
- Batched Processing
- Streaming Processing

# Today's topic: Parallelism

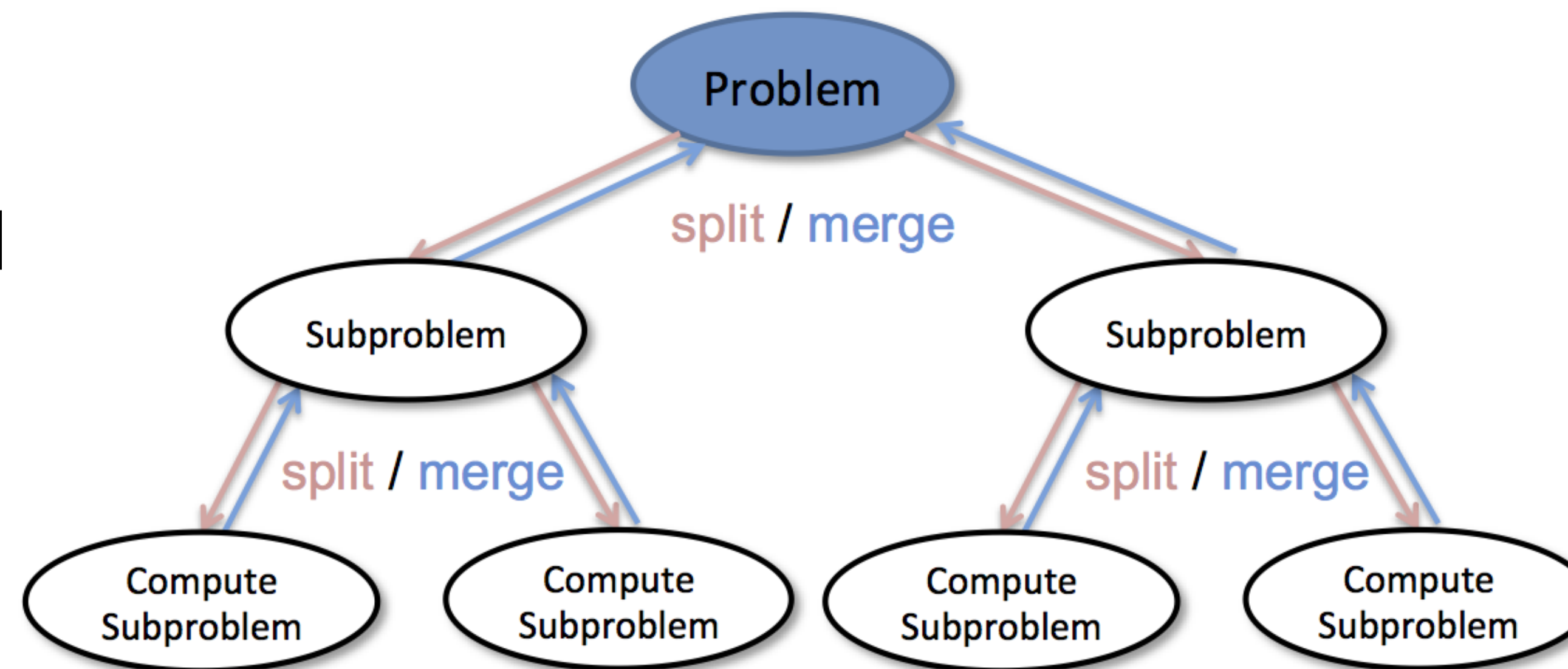
- Express data processing in abstraction
- Parallelisms

# Parallel Data Processing

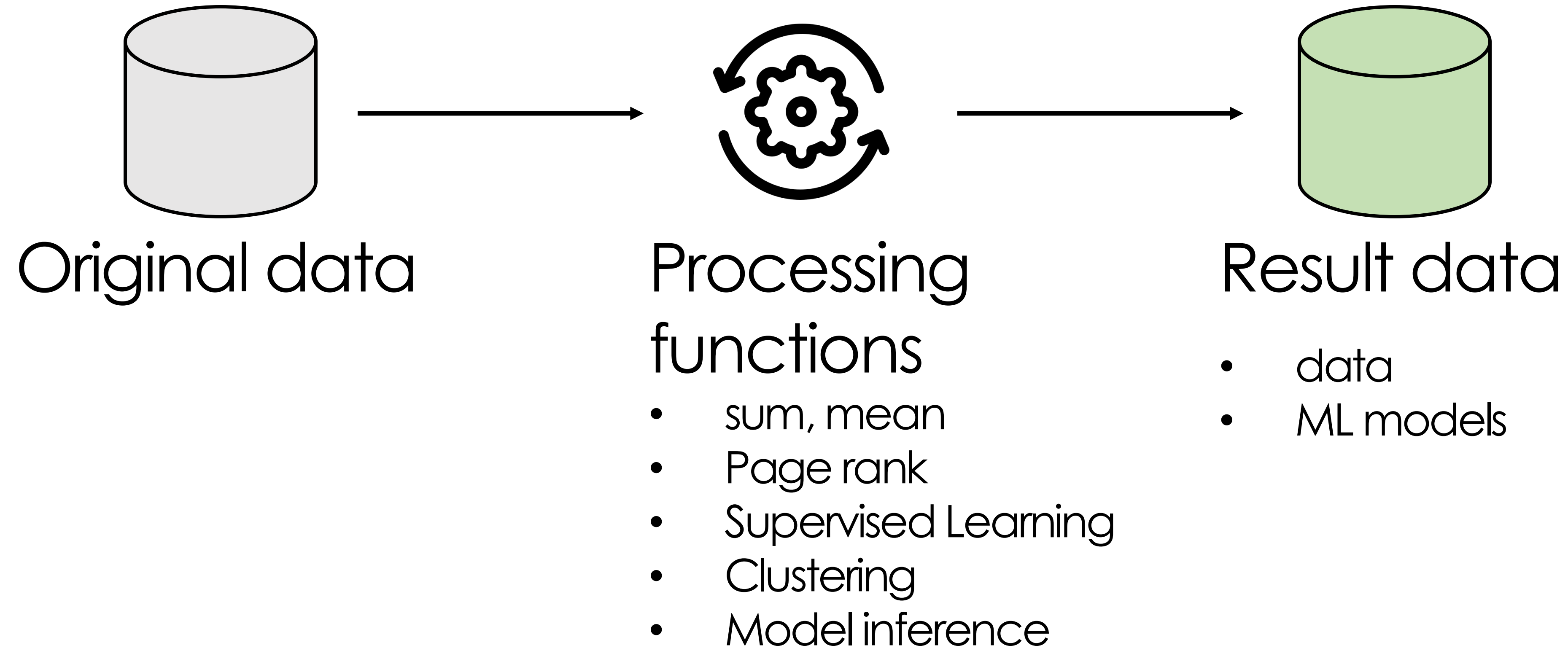
**Central Issue:** Workload takes too long for one processor!

**Basic Idea:** Split up workload across processors and perhaps also across machines/workers (aka “Divide and Conquer”)

Remind you of PA1  
(hope you've  
enjoyed it)



# Data Processing: Abstraction



Q: How to represent various processing functions?

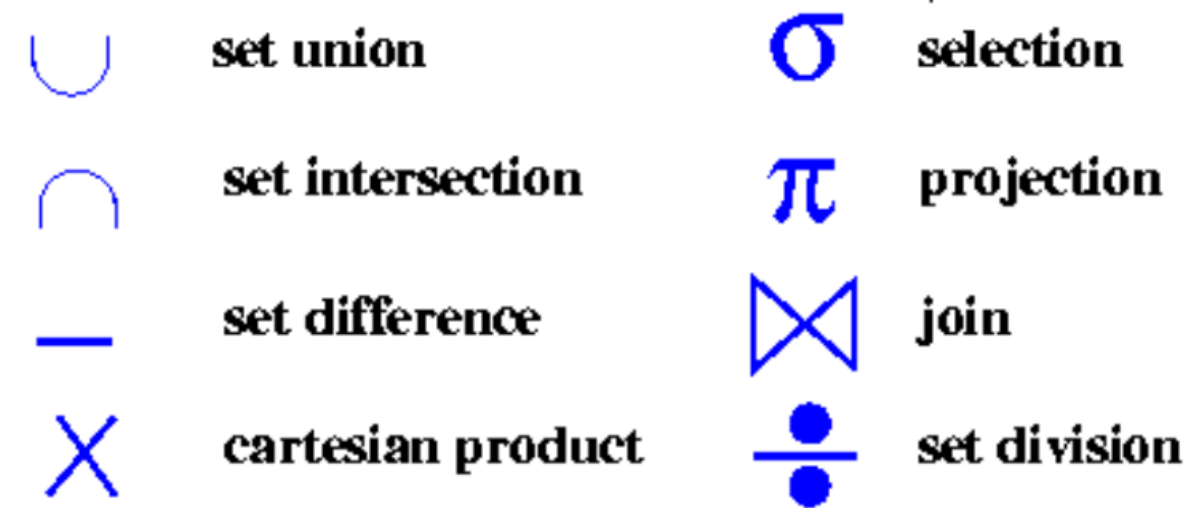


# How to Express Arbitrarily Complex Processing Functions?

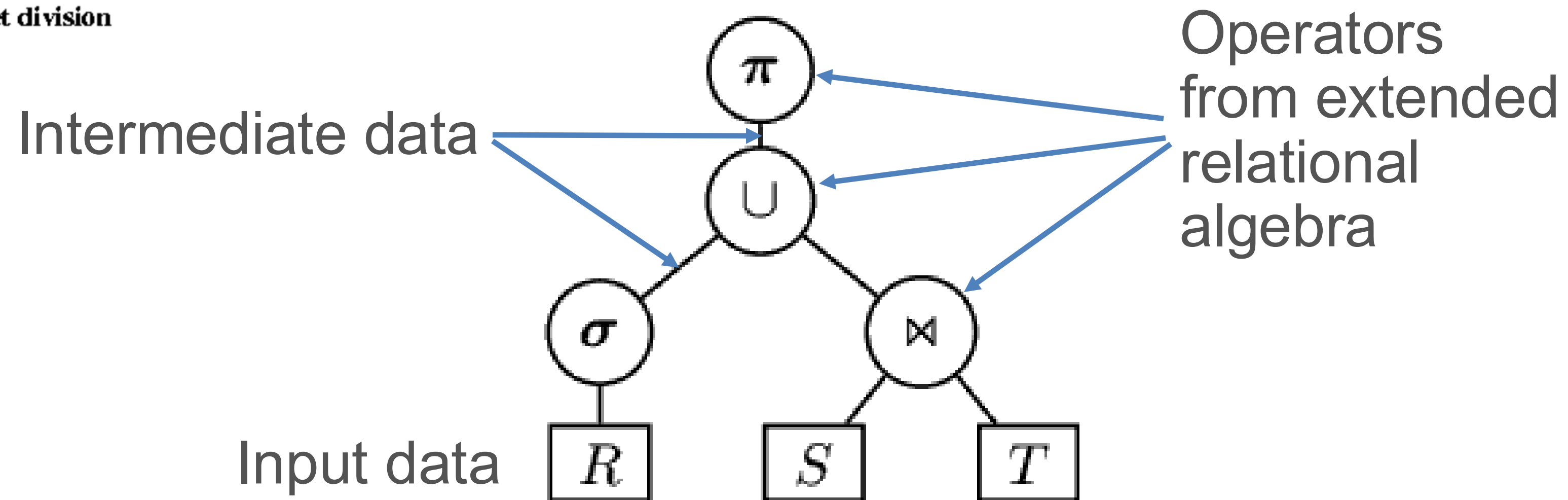
**Dataflow Graph:** common in parallel data processing

- A **directed** graph representation of a program
  - **Vertices:** abstract operations from a restricted set of computational primitives:
  - **Edges:** data flowing directions (hence data dependency)
- Examples
  - Relational dataflows: RDBMS, Pandas, Modin
  - Matrix/tensor dataflows: NumPy, PyTorch, TensorFlow
- Enables us to reason about data-intensive programs at a higher level

# Example: Relational Dataflow Graph



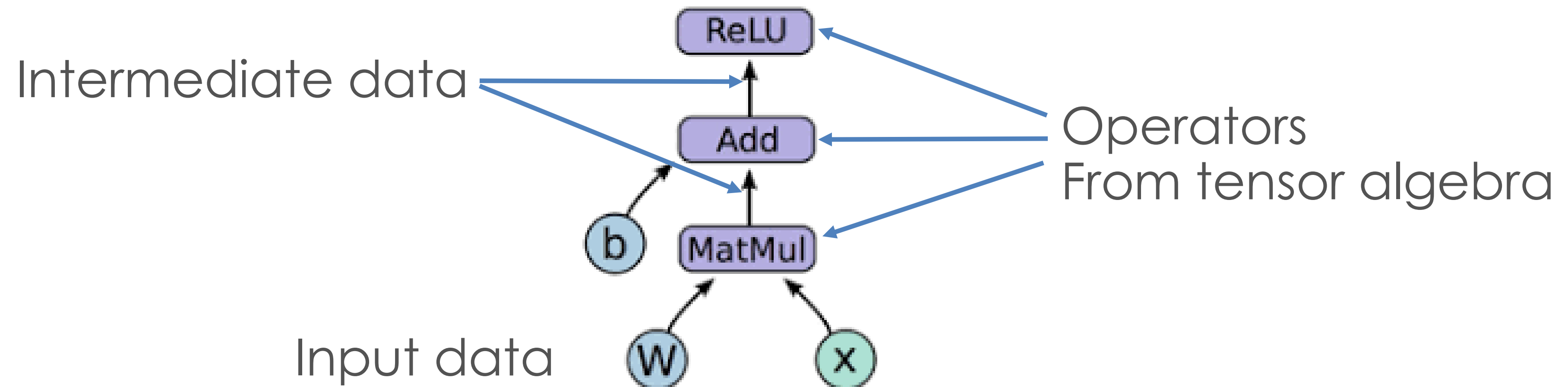
$$\pi(\sigma(R) \cup S \bowtie T)$$



Aka **Logical Query Plan** in the DB systems world

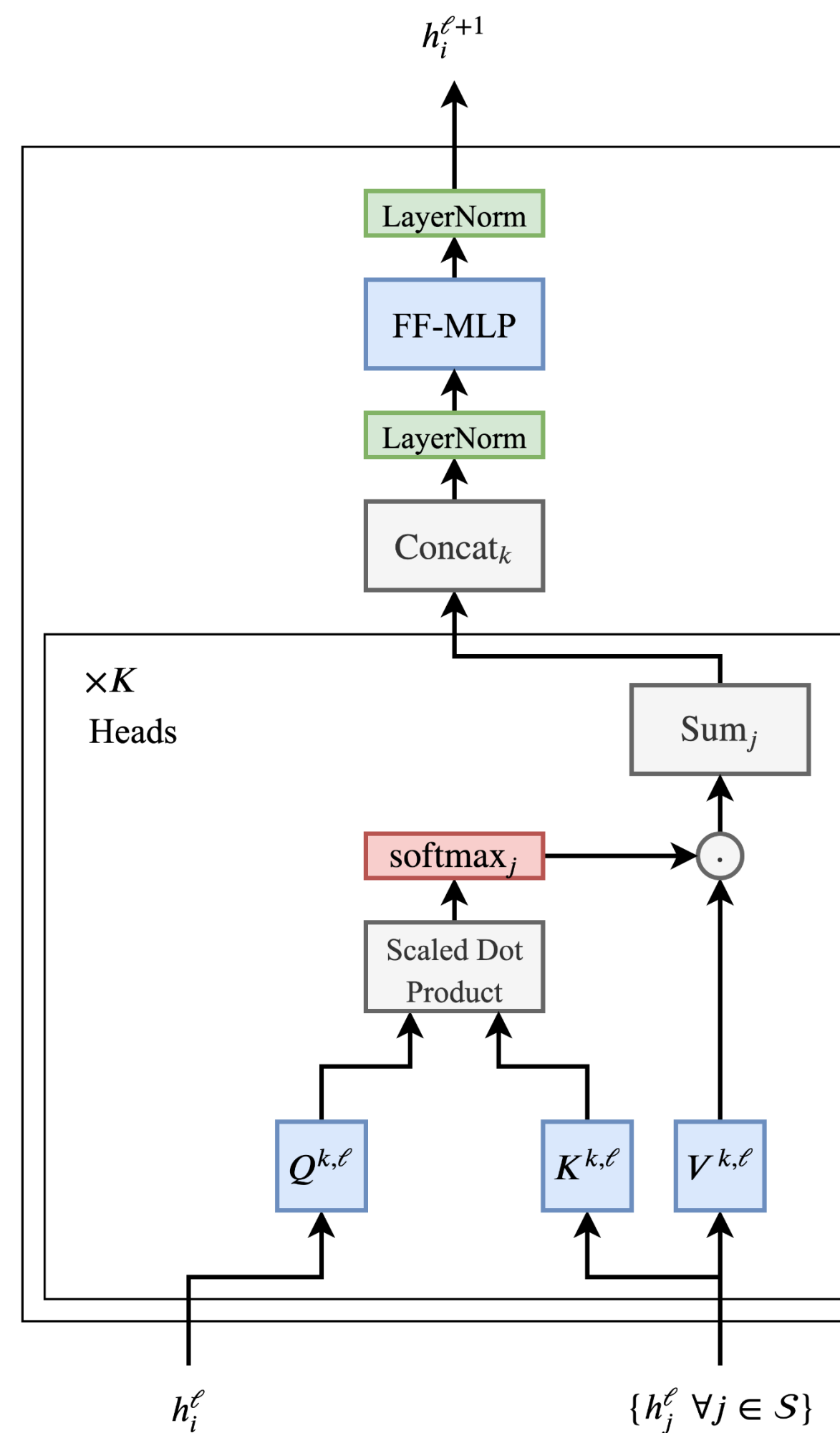
# Example: Machine Learning Dataflow Graph

$$\text{ReLU}(WX + b)$$



Aka **Neural network computational graph** in ML systems

# What is ChatGPT's dataflow graph Looking like?



# Parallelism

**Central Issue:** Workload takes too long for one processor!

**Basic Idea:** Split up workload across processors and perhaps also across machines/workers (aka “Divide and Conquer”)

**Key parallelism paradigms in data systems**

- assuming there will be coordination:

func	data		
	Shared	Replicated	Partitioned
Replicated	N/A (rare cases)		Data parallelism
Partitioned	Task parallelism		Hybrid parallelism