

Misc

Misc

- $\cos(a+b)+i\sin(a+b) = e^{i(a+b)} = [\cos(a)+i\sin(a)][\cos(b)+i\sin(b)] = \cos(a)\cos(b) - \sin(a)\sin(b) + i[\cos(a)\sin(b) + \cos(b)\sin(a)]$
 - $\cos(2\theta) = \cos^2(\theta) - \sin^2(\theta)$
 - $\sin(2\theta) = 2\cos(\theta)\sin(\theta)$
- $\sin(\pi-x) \equiv \sin(x)$
- $\cos(-x) \equiv \cos(x)$
- $\sin(x\pm\pi) \equiv -\sin(x)$
- $\cos(x\pm\pi) \equiv -\cos(x)$
- $\tan(x\pm\pi) \equiv \tan(x)$
- Something that was inside a trig function (eg $2x$ in $\sin(2x)$) should not be simplified further until all solutions have been found — adjust the range instead

Number systems

Bases

- Base conversion: Divide by the new base, record the remainder, then repeat with the quotient, once you have a quotient of 0 stop and write the number left-to-right using the remainders bottom-to-top
- If a divides b , then to convert between base a and b it suffices to convert each digit to/from b and concatenate the results

Euclid's division algorithm

```
divide(a, b)
    r := a
    q := 0
    while r >= b
        r -= b
        q++
    //a = qb + r (a/b gives quotient q and remainder r)
    return (q, r)
```

Greatest common divisor

gcd(m, n)

```
assert m >= 0 && n >= 0
// Or just assert ordering
if m > n
    m, n = n, m
if m == 0
    return n
return gcd(n mod m, m)
```

- This works because if d divides $n=qm+r$ then d divides qm and r

Modular arithmetic

- a and b are congruent mod $n \Leftrightarrow \exists k \in \mathbb{Z}: a=b+kn \Leftrightarrow n$ divides $a-b$
- **If $x \bmod n = a$ and $y \bmod n = b$:**
 - $x+y \bmod n = a + b \bmod n$
 - $xy \bmod n = ab \bmod n$
 - $x^k \bmod n = a^k \bmod n$
 - $x-y \bmod n = a - b \bmod n$
 - **Division is complicated**
 - $z^{xy} \bmod n = (z^x \bmod n)^y \bmod n$
- *Fermat's little theorem: For all integers n and p , if p is prime and n and p are coprime, then $n^{p-1} \bmod p = 1$*
- Interpreting an n -bit two's complement number x as unsigned gives $x \bmod 2^N$ — the above is why normal arithmetic works with two's complement

Sets of numbers

Subsets of the reals

- $x \in \mathbb{Q} \Leftrightarrow \exists m \in \mathbb{Z}: \exists n \in \mathbb{N} \setminus \{0\}: z = m/n$, and m and n are coprime
- x is algebraic iff x is real and there exists a finite polynomial with rational (or equivalently integer) coefficients with a root x
- x is transcendental iff it is real and it is not algebraic

The real numbers: Axioms

- Commutativity of addition and multiplication
- Associativity of addition and multiplication
- Distributivity of multiplication over addition
- Existence of additive identity 0
- Existence of multiplicative identity 1
- 0 and 1 are distinct
- All the above hold for the naturals also
- Existence of additive inverse: $\forall x. \exists -x: x+(-x) = 0$
- All the above hold for the integers also
- Existence of multiplicative inverse: $\forall x. \exists x^{-1}: x \cdot (x^{-1}) = 1$
- All the above hold for the rationals and algebraics also
- Transitivity: $\forall x, y, z. (x < y \wedge y < z) \Rightarrow x < z$
- Trichotomy: **For all x, y , exactly one of these is true: $x < y, y < x, x = y$**
- Preservation of ordering under addition: $\forall x, y, z. x < y \Rightarrow x + z < y + z$
- Preservation of ordering under multiplication: $\forall x, y, z. (x < y \wedge 0 < z) \Rightarrow x \cdot z < y \cdot z$
- **Completeness: Every non-empty subset that has an upper bound has a least upper bound in the reals**

The real numbers: Completeness

- Let $S \subseteq \mathbb{R}$, $S \neq \emptyset$, and \leq be the standard total ordering over \mathbb{R}
 - $u \in \mathbb{R}$ is an upper bound of S iff $\forall x \in S. x \leq u$
 - S is bounded iff there exists an upper bound of S
 - $u \in \mathbb{R}$ is the least upper bound (supremum) of S iff it is the least element of the set of upper bounds of S (for all upper bounds u' , $u \leq u'$)

- *There are dual notions of lower bound, greatest lower bound, and infimum*

- **Theorem (Archimedean property of the reals):** $\forall a, b \in \mathbb{R}_{>0}. \exists n \in \mathbb{N}: na > b$

Proof: Assume for the sake of contradiction that there exists positive reals a, b such that for all natural n $na \leq b$

Then, b is an upper bound of $\{na: n \in \mathbb{N}\}$

By the completeness of the reals, there exists a supremum l of $\{na: n \in \mathbb{N}\}$

Hence, $\forall n \in \mathbb{N}. na \leq l$

Hence, $\forall n \in \mathbb{N}. (n+1)a \leq l$. So, $\forall n \in \mathbb{N}. na \leq l-a$

Hence, $l-a$ is an upper bound. As a is positive, $l-a < l$. Hence, there exists an upper bound smaller than the supremum?!

Complex numbers

- $(a+bi)(c+di)=ac-bd+(ad+bc)i$
- If $z=a+bi$, z^* (can also be denoted by an overbar) $=a-bi$
 - $(z+w)^* = z^* + w^*$
 - $(zw)^* = z^*w^*$
 - $(z/w)^* = z^*/w^*$
- If $z=re^{i\theta}$, then $z=r[\cos(\theta) + i\sin(\theta)]$ and $z^*=re^{-i\theta}$
 - If $z=x+iy$, $r^2=x^2+y^2$ and $\tan(\theta)=y/x$
 - θ is measured anticlockwise from the real axis — a diagram is necessary to establish the relationship between $\arctan(y/x)$ and θ
- To evaluate z_1/z_2 multiply top and bottom by z_2^* to obtain $z_1/z_2=(z_1z_2^*)/((|z_2|)^2)$
- $|z_1z_2|=|z_1||z_2|$
- Triangle inequality: $|z_1+z_2| \leq |z_1| + |z_2|$
- $||z| - |w|| \leq |z-w|$
- **Fundamental theorem of algebra: Every polynomial of degree n has exactly n (not necessarily distinct) roots in \mathbb{C}**

Vectors

Vectors

- Vectors u and v are equal iff their components are equal iff they have the same length and are parallel and in the same direction
- If a and b are position vectors, then $c=a+b$ is the position vector s.t. $OACB$ is a parallelogram or equivalently is the point obtained by top and tailing a and b
- **Pythagoras extends to higher dimensions — $|(a_1, \dots, a_n)|^2 = \text{sum of } (a_i^2)$**
- **$(a_1, \dots, a_n) \cdot (b_1, \dots, b_n) = \text{sum of } (a_i b_i) = |a||b|\cos(\theta)$ where θ is angle between a and b**

Proof for $n=2$: Let α be the angle made between a and the origin and β be the angle made between b and the origin. Without loss of generality, assume $\alpha \leq \beta$. Then, $\theta = \beta - \alpha$.

Thus, $\cos(\theta) = \cos(\beta)\cos(\alpha) + \sin(\beta)\sin(\alpha) = (b_1/|b|)(a_1/|a|) + (b_2/|b|)(a_2/|a|) = a_1 b_1 / |a||b| + a_2 b_2 / |a||b|$.

So, $a_1 b_1 + a_2 b_2 = |a||b|\cos(\theta)$

- a and b are orthogonal (perpendicular) iff $a \cdot b = 0$
- $|\lambda v| = |\lambda||v|$

Subspaces

- $\text{span}(\{u_1, \dots, u_m\}) = \{a_1 u_1 + \dots + a_m u_m : a_1, \dots, a_m \in \mathbb{R}\} = \text{the set of all linear combinations of } u_1, \dots, u_m$
 - Equivalently $\text{span}(\{u_1, \dots, u_m\}) = \text{the vector } a_1 u_1 + \dots + a_m u_m \text{ where } a_1, \dots, a_m \text{ are real parameters}$
- A finite subset S of \mathbb{R}^n is a subspace of \mathbb{R}^n iff S is: non-empty, closed under addition ($u, v \in S \Rightarrow u+v \in S$), and closed under scalar multiplication ($u \in S$ and $\lambda \in \mathbb{R} \Rightarrow \lambda u \in S$)
 - Equivalently, S is a subspace of \mathbb{R}^n iff $S \subseteq \mathbb{R}^n$ and $S = \text{span}(S) \neq \emptyset$
 - Hence, if S is a non-empty finite subset of \mathbb{R}^n , then $\text{span}(S)$ is a subspace of \mathbb{R}^n
 - Due to closure under scalar multiplication, every subspace of \mathbb{R}^n contains the zero vector

Linear independence

- $\{u_1, \dots, u_m\}$ is **linearly dependent** iff a vector can be written as a linear combination of the others **iff there is a linear combination with coefficients not all zero that gives the zero vector**
- $\{u_1, \dots, u_m\}$ is **linearly independent** iff it is not linear dependent **iff (if a linear combination gives the zero vector then the coefficients are all zero)**

Bases

- A set of vectors U is a basis of a subspace S iff U is linearly independent and $\text{span}(U)=S$
- If $V=\{v_1, \dots, v_m\}$ spans S , then any linearly independent subset of S contains at most m vectors
- Any two bases of the same subspace have the same number of elements
Proof: Let $U=\{u_1, \dots, u_m\}$ and $W=\{w_1, \dots, w_p\}$ be bases for S
As U spans S and W is a linearly independent subset of S , $m \leq p$
As W spans S and U is a linearly independent subset of S , $p \leq m$
Thus, $p=m$
- Dimension of a subspace = cardinality of any basis for the subspace
- If $U \subseteq S$ and $|U| > \dim(S)$, then U is linearly dependent

Matrices

Misc

- A matrix of order $m \times n$ has m rows and n columns
- a_{ij} is the element in the i th row and j th column
- A matrix is diagonal iff every element not on the leading diagonal is zero (elements on leading diagonal may also be zero)
 - Matrix multiplication of diagonal matrices reduces to element wise multiplication
 - *Determinant of a diagonal matrix = product of elements along diagonal*
- Matrix multiplication is associative
- Matrix multiplication distributes over matrix addition
- A^T is the transpose of A — the rows of A^T are the columns of A and the columns of A^T are the rows of A
 - $(AB)^T \equiv B^T A^T$
 - $(A+B)^T \equiv A^T + B^T \equiv B^T + A^T$ as matrix addition is commutative
- If $C=AB$, the element in row i and column j of C is the dot product of row i of A and column j of B — row, row, row your boat down the column fall
- Can either pre-multiply a matrix to both sides or post-multiply a matrix to both sides — can not place it wherever is convenient
- $(ABC\dots)^{-1} = \dots C^{-1} B^{-1} A^{-1}$

Row operations

- **The elementary row operations are: Swap two rows, multiply a row by a non-zero real, add a non-zero real multiple of a row to another**
- $A \sim B$ (A and B are row equivalent) iff A can be transformed to B using a series of elementary row operations
- Each elementary row operation can be written as pre-multiplication by a(n invertible) matrix
- A matrix is in row-echelon form iff the first non-zero entry in each row is further right than the first non-zero entry in the previous row
- Creating an augmented matrix by including the RHS then reducing to row echelon form simplifies a system of simultaneous equations

Row operations and the determinant

- **Swapping two rows causes the sign of the determinant to flip**

Proof sketch: Induction on n . Base case $n=2$. Inductive step, calculate determinant using a row that wasn't in the swap

- **Add a multiple of a row to another has no impact on the determinant**

Proof sketch: Corresponding matrix can fairly easily be shown to have $\det = 1$ (as **a matrix with non-zero elements only on the diagonal has $\det = \text{product of diagonal elements}$**)

- Corollary: A row is a multiple of another $\Rightarrow \det = 0$

- **Multiplying a row (or a column) by a constant k increases the determinant by a factor of k**

Proof: Trivial

Matrix inverse: Row operations

- **Theorem: If a series of elementary row operations transforms a (square) matrix A into I , then A is invertible and the same sequence transforms I into A^{-1}**

Proof: We have that $E_n \dots E_1 A = I$. Let $E = E_n \dots E_1$

As E is a product of invertible matrices, it is invertible

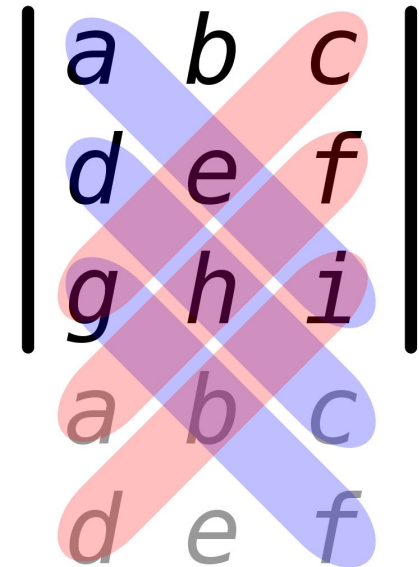
$$AE = IAE = E^{-1}EAE = E^{-1}IE = I$$

So, $A^{-1}AE = A^{-1}I$ and thus $E = A^{-1}$

- Augmenting a matrix by including the identity matrix then carrying out row operations to turn it into the identity matrix causes the augmented part to become the inverse

Determinant of a 3x3 matrix: The rule of Sarrus

- Only works for 3x3
1. Write out the first 2 rows again below the matrix
 2. Working your way down the left hand column, sum the products of the (first 3 elements of the) leading diagonals
 3. Working your way down the right hand column, sum the products of the (first 3 elements of the) trailing diagonals
 4. Determinant = sum of leading diagonals – sum of trailing diagonals



Matrix inverse: Cofactors

$$\begin{bmatrix} + & - & + \\ - & + & - \\ + & - & + \end{bmatrix}$$

Multiply the minor by +1 or -1 accordingly to find the cofactor

1. Create a new matrix C where each element is the cofactor of that element in the original matrix
2. Write down C^T
3. Write down the determinant of the original matrix
4. Inverse = $1/|\text{original}| \times C^T$

Linear independence of vectors

- **A set of n vectors is linearly independent iff the matrix it is the column vectors of has non-zero determinant**

Proof sketch: Linearly independent iff $Ua = 0$ doesn't have any solutions other than $a=0$ where U is the matrix, a is the vector of coefficients, and 0 is the zero vector

Only solution is $a=0$ iff U is invertible iff U has non-zero determinant

Linear transformations and coordinates

- A function $T: \mathbb{R}^m \mapsto \mathbb{R}^n$ is a linear transform iff it preserves addition ($T(u+v) \equiv T(u)+T(v)$) and preserves scalar multiplication ($T(\lambda u) \equiv \lambda T(u)$)
 - Hence, must map zero vector to itself
- **Theorem: A function is a linear transform iff it can be written as multiplication by a matrix**

Proof sketch: Matrix multiplication distributes over addition and preserves scalar multiplication
- **Coordinates of a vector v with respect to a basis B = vector of coefficients of linear combination of B that gives v**
- **To construct a matrix from a linear transform T : Pick bases for \mathbb{R}^m and \mathbb{R}^n $V=\{v_1, \dots, v_m\}$ and $W=\{w_1, \dots, w_n\}$ and let the i^{th} column be the coordinates of $T(v_i)$ with respect to W**
- **A transition matrix (and its inverse) allows conversion between coordinates in two bases for the same subspace**
- **Transition matrix = matrix for linear transform using identity function**

Eigenvectors

- \underline{v} is an eigenvector of $M \Leftrightarrow M\underline{v} = \lambda\underline{v}$ where λ is a scalar constant (the eigenvalue)
 - $M\underline{v} = \lambda\underline{v} \Leftrightarrow IM\underline{v} = I\lambda\underline{v} \Leftrightarrow IM\underline{v} - I\lambda\underline{v} = 0 \Leftrightarrow (M - I\lambda)\underline{v} = 0$
- **To find eigenvalues: Solve $\det(M - \lambda I) = 0$ for λ**
- **To find eigenvector sub in λ and solve $M\underline{v} = \lambda\underline{v}$ for \underline{v}**
- **A root in the eigenvalue equation that is repeated n times could have up to n non-parallel eigenvectors associated with it**
- **If there are non-real solutions, include them**
- *Sum of the leading diagonal of a matrix \equiv sum of its eigenvalues — this is a useful check*
- *Product of eigenvalues $\equiv \det(M)$ — this is a useful check*

Diagonalisation

- An $n \times n$ matrix is diagonalisable \Leftrightarrow it has n non-parallel eigenvectors
- **Let V be the matrix where each column is an eigenvector of M .
Let D the diagonal matrix of eigenvalues for the eigenvectors in the corresponding columns in U .
Then, $V^{-1}MV=D$**
 - **Corollary: $VDV^{-1}=M$ and thus $M^k=VD^kV^{-1}$**

Sequences

Convergent sequences

- Sequence a_n converges to l ($a_n \rightarrow l$) iff $\forall \varepsilon \in \mathbb{R}_{>0}. \exists N \in \mathbb{N}: \forall n \in \mathbb{N}_{>N}. |a_n - l| < \varepsilon$
 - To prove: Construct a suitable N in terms of **epsilon** and show this provides a sufficient condition
- If a_n converges to l , then every subsequence of a_n also converges to l
 - Contrapositive: If there exists a subsequence of a_n that does not converge to l , then a_n does not converge to l
 - Corollary: If there exists a pair of subsequences that converge to different limits, then a_n is divergent
 - Subsequence = a sequence (of infinitely many terms) obtained by removing some terms without adding any
- Squeeze rule: If a_n and b_n both converge to the same limit l and $\exists N \in \mathbb{N}: \forall n \in \mathbb{N}_{>N}. a_n \leq c_n \leq b_n$, then c_n also converges to l
- If a_n, b_n, c_n converge to A, B, C respectively:
 - $|a_n| \rightarrow A$
 - $\lambda a_n \rightarrow \lambda A$
 - $a_n + b_n \rightarrow A + B$
 - $a_n b_n \rightarrow AB$
 - $1/a_n \rightarrow 1/A$

Properties of sequences

- a_n **diverges** (to (positive) infinity) **iff** $\forall K \in \mathbb{R}. \exists N \in \mathbb{N}: \forall n \in \mathbb{N}_{>N}. a_n > K$
- a_n is decreasing **iff** $\forall n \in \mathbb{N}. a_n \geq a_{n+1}$
- a_n is increasing **iff** $\forall n \in \mathbb{N}. a_n \leq a_{n+1}$
- a_n is **bounded above** **iff** $\exists U \in \mathbb{R}. \forall n \in \mathbb{N}. a_n \leq U$
- a_n is **bounded below** **iff** $\exists L \in \mathbb{R}. \forall n \in \mathbb{N}. L \leq a_n$
- a_n is **bounded** **iff** it is bounded above and below
- If a_n is convergent, then it is bounded
- If a_n is increasing and bounded above, then it is convergent
- If a_n is decreasing and bounded below, then it is convergent
- a_n **oscillates** **iff** it is neither convergent nor divergent

Asymptotics

- $a_n \in O(b_n)$ iff $\exists c, N \in \mathbb{R}: \forall n \in \mathbb{N}_{>N} [|a_n| \leq c|b_n|]$ — b_n is an asymptotic upper bound for a_n
- $a_n \in \Omega(b_n)$ iff $\exists c, N \in \mathbb{R}: \forall n \in \mathbb{N}_{>N} [|a_n| \geq c|b_n|]$ — b_n is an asymptotic lower bound for a_n
- $a_n \in \Theta(b_n)$ iff $a_n \in O(b_n)$ and $a_n \in \Omega(b_n)$
- **A useful trick:** $[f(n)]^k + [f(n)]^{k-1} + \dots + [f(n)]^0 \leq [f(n)]^k + \dots + [f(n)]^k$ if $f(n) \geq 1$
E.g. $n^2+2n+1 \in O(n^2)$ as $n^2+2n+1 \leq n^2+2n^2+n^2$ (for $n>1$) $= 4n^2$

Recurrence relations

- The homogeneous second-order linear recurrence $ax_n + bx_{n-1} + cx_{n-2} = 0$ has auxiliary equation $a\lambda^2 + b\lambda + c = 0$
 - **If auxiliary equation has distinct roots r_1, r_2 : $x_n = Ar_1^n + Br_2^n$**
 - **If auxiliary equation has repeated root r : $x_n = Ar^n + Bnr^n$**
 - A and B are found using initial conditions
- The homogeneous second-order linear recurrence $ax_n + bx_{n-1} + cx_{n-2} = f(n)$ is solved by solving the homogeneous case and adding a particular solution **$p(n)$ s.t. $ap(n) + bp(n-1) + cp(n-2) \equiv f(n)$**
 - If $f(n)$ is a k th degree polynomial, use a general k th degree of polynomial in n as $p(n)$
 - *Multiply $p(n)$ through by n if a root of auxiliary equation is 1 — multiply through by n^2 if 1 is a repeated root*

Series

Series

- sum from $i=0$ to $i=\infty$ of a_i converges to t iff the sequence of partial sums $s_n = \text{sum from } i=0 \text{ to } i=n \text{ of } a_i$ converges to t
- a_n converges to 0 is necessary but not sufficient for $\sum a_n$ converges
- $\sum |a_n|$ is sufficient but not necessary for $\sum a_n$ converges
- If $\sum a_n$ and $\sum b_n$ converge to S and T :
 - $\lambda \sum a_n \rightarrow \lambda S$
 - $\sum a_n + b_n \rightarrow S + T$
- Comparison test: If $\sum b_n$ converges and $\exists N \in \mathbb{N}: \forall n \in \mathbb{N}_{>N}. 0 \leq a_n \leq b_n$, then a_n also converges
 - Contrapositive: If $\sum a_n$ diverges and $\exists N \in \mathbb{N}: \forall n \in \mathbb{N}_{>N}. 0 \leq a_n \leq b_n$, then b_n also diverges
 - It is not a coincidence that this looks a lot like the squeeze rule
- Ratio test: Let $|a_{n+1}/a_n| \rightarrow k$, then:
 - $0 \leq k < 1 \Rightarrow \sum a_n$ converges
 - $k > 1$ (including ratio diverges) $\Rightarrow \sum a_n$ diverges
 - $k=1$ \Rightarrow was a waste of time, can't tell anything

Power series

- A power series is a series of the form $\sum a_n x^n$
- Theorem: If $\exists R \in \mathbb{R}_{\geq 0} : \sum a_n R^n$ converges then, $\forall x \in \mathbb{R} [|x| < R \Rightarrow \sum a_n x^n \text{ converges}]$
 - The largest such R is called the radius of convergence — for $|x| < R$ $\sum a_n x^n$ converges and for $|x| > R$ $\sum a_n x^n$ diverges — behaviour at $x = \pm R$ is indeterminate
 - **Ratio test is a good way of finding**
- Let $f(x) = \sum a_n x^n$ and $g(x) = \sum b_n x^n$ have domains $|x| < R_1$ and $|x| < R_2$ where R_1 and R_2 are the radii of convergence. Let $R = \min(R_1, R_2)$. Then for all $|x| < R$:
 - If $f(x) = g(x)$, then $a_n = b_n$ for all n
 - **$f(x) + g(x) = \sum (a_n + b_n) x^n$**
 - $\lambda f(x) = \sum \lambda a_n x^n$
 - **$f(x)g(x) = \sum (a_0 b_n + a_1 b_{n-1} + \dots + a_n b_0) x^n$**
 - These are useful when considering Taylor series
- **Binomial theorem: $\sum \binom{q}{n} x^n$ converges to $(1+x)^q$ with radius of convergence $|x| < 1$**

Long division

- Let $s/t \in (0; 1)$:
 - $r_0 = 10s$
 - $q_{i+1} = r_i \text{ div } t$
 - $r_{i+1} = 10(r_i \bmod t)$
 - $s/t = \sum q_i / 10^i$ starting at $i=1$ up to $i=\text{inf}$ (or equivalently up to i s.t. $r_i=0$)
 - For a finite approximation: $s/t = q_1/10^1 + \dots + q_n/10^n (+ r_n/t10^{n+1})$

Calculus

Limits

Limits

- Let f be a function with domain of I or $I \setminus \{a\}$
 - $\lim_{x \rightarrow a^-} f(x) = l$ iff for every sequence x_n in I s.t. $x_n \rightarrow a$ and x_n is bounded above by a , the sequence $f(x_n) \rightarrow l$
 - $\lim_{x \rightarrow a^+} f(x) = l$ iff for every sequence x_n in I s.t. $x_n \rightarrow a$ and x_n is bounded below by a , the sequence $f(x_n) \rightarrow l$
 - $\lim_{x \rightarrow a} f(x) = l$ iff for every sequence x_n in I s.t. $x_n \rightarrow a$ and x_n does not contain a , the sequence $f(x_n) \rightarrow l$
 - $\lim_{x \rightarrow a} f(x) = l$ iff $\lim_{x \rightarrow a^-} f(x) = l$ and $\lim_{x \rightarrow a^+} f(x) = l$
- Same rules apply as for convergent sequences

Continuity

- **f is continuous at $a \in I$ iff $\lim_{x \rightarrow a} f(x) = f(a)$**
- **f is continuous iff for all $a \in I$ $\lim_{x \rightarrow a} f(x) = f(a)$**
- If f is continuous at a and g is continuous at $f(a)$, then $g \circ f$ is continuous at a
- **Intermediate value theorem: If f is continuous and is defined at a and b and $f(a)$ and $f(b)$ have opposite signs, then $f(c)=0$ for some $c \in (a; b)$**
- **Extreme value theorem: If f is defined for and continuous over $[a; b]$, then $\exists m, M \in [a; b]: \forall x \in [a; b]. f(m) \leq f(x) \leq f(M)$ — all intervals are closed**

Differentiation

Differentiation

- $f'(a) = \lim_{h \rightarrow 0} (f(a+h)-f(a))/h = \lim_{x \rightarrow a} (f(x)-f(a))/(x-a)$
 - **f is differentiable at a iff this limit exists**
 - **f is differentiable iff it is differentiable at every point in its domain**
 - **f is continuous at a is necessary but not sufficient for f is differentiable at a**
- Proof that f is differentiable at a \rightarrow f is continuous at a: We have that $\lim_{x \rightarrow a} (f(x)-f(a))/(x-a)$ exists, let it be c.

$$\text{Then, } \lim_{x \rightarrow a} f(x)-f(a)=c \cdot \lim_{x \rightarrow a} x-a = 0$$

So, $\lim_{x \rightarrow a} f(x)-\lim_{x \rightarrow a} f(a)=0$ and thus $\lim_{x \rightarrow a} f(x)=f(a)$ as f(a) is a constant

- If f' and g' exist then:
 - $(f+g)'=f'+g'$
 - $(\lambda f)'=\lambda f'$
 - $(f \cdot g)'=f' \cdot g+g' \cdot f$
 - $(g \circ f)'=(g' \circ f) \cdot f'$
 - $(f/g)'=(f' \cdot g-f \cdot g')/g^2$
- In partial differentiation (δ instead of d), all variables not on the bottom of the derivative operator are treated as constants

Using differentiation

- a is a turning point iff there exists a closed interval I near a s.t. $\forall x \in I. f(x) \leq f(a)$ or $\forall x \in I. f(a) \leq f(x)$ iff a is a local minimum or a local maximum
- Global maximum and minimum may be endpoints instead of turning points
- a is a stationary point iff $f'(a)=0$ iff f is either a turning point or a point of inflection
- n th derivative test: If $f'(a)=\dots=f^{(n-1)}(a)=0$ and $f^{(n)}$ is continuous then:
 - **n is even and $f^{(n)}(a) > 0 \Rightarrow a$ is a local minimum**
 - **n is even and $f^{(n)}(a) < 0 \Rightarrow a$ is a local maximum**
 - **n is odd and $f^{(n)}(a) \neq 0 \Rightarrow a$ is a point of inflection**
- Turning point theorem: If f is differentiable and has a turning point at a , then $f'(a)=0$
- Rolle's theorem: **If f is continuous over $[a;b]$ and differentiable over $(a;b)$ and $f(a)=f(b)$, then $\exists c \in (a; b): f'(c)=0$**

Proof:

By extreme value theorem f has extreme values in $[a;b]$

Case f is a constant function over $(a;b)$: Then, $f'(c)=0$ for any choice of c

Case f is not a constant function over $(a;b)$: Then, extreme values must be distinct, thus $\exists c \in (a; b)$ s.t. c is a turning point and so by turning point theorem $f'(c)=0$

- Mean value theorem: **If f is continuous over $[a;b]$ and differentiable over $(a;b)$, then $\exists c \in (a; b): f'(c)=(f(b)-f(a))/(b-a)$**

Proof: Consider $h(x)=f(x)-(f(b)-f(a)) \cdot (x-a)/(b-a)$

Deduce that $h(a)=f(a)$ and $h(b)=f(a)$ ($=h(a)$). Hence, by Rolle's theorem, $\exists c \in (a; b): h'(c)=0$.

Thus, as $h'(c)=f'(c)-(f(b)-f(a))/(b-a)$, $f'(c)=(f(b)-f(a))/(b-a)$

Sketching rational functions

- For $y=(ax+b)/(cx+d)$ and $y=(ax^2+bx+c)/(dx^2+ex+f)$:
 - x-intercept \Leftrightarrow numerator=0
 - y-intercept $\Leftrightarrow x=0$
 - Vertical asymptote \Leftrightarrow denominator=0
 - Horizontal asymptote $\Leftrightarrow x \rightarrow \infty \Leftrightarrow y$ is approaching a/c in first case and a/d in the second case
- To find (y-values of) stationary points set the equation equal to k (to find intersection points with $y=k$) then find discriminant, set that equal to 0, and solve for k
 - The set of values that y cannot take is described by discriminant < 0
- Oblique asymptotes: If $y=(ax^2+bx+c)/(ex+f)$:
 - All above rules apply except for the one for horizontal asymptotes
 - Instead, y tends towards $(ax^2+bx+c)/(ex+f)$ i.e. polynomial division but discard the remainder — this process works in all cases dividing coefficients is a shortcut that only works if $f=0$

L'hospital's rule

- Weak L'hospital's rule: **If $f(a)=g(a)=0$ and $g'(a)\neq 0$, then**

$$\lim_{x \rightarrow a} f(x)/g(x) = f'(a)/g'(a)$$

Proof: $\lim_{x \rightarrow a} f(x)/g(x) = \lim_{x \rightarrow a} (f(x)-f(a))/(g(x)-g(a))$
 $= (f'(a)/g'(a)) \lim_{x \rightarrow a} (x-a)/(x-a)$ as limits preserve multiplication
 $= f'(a)/g'(a)$

- L'hospital's rule: **If $f(a)=g(a)=0$ and f and g are differentiable, then**

$$\lim_{x \rightarrow a} f(x)/g(x) = \lim_{x \rightarrow a} f'(x)/g'(x)$$

- *Strong l'hospital's rule: If f and g are differentiable and $f(a)=g(a)=0$ or $f(a)=g(a)=\pm\infty$, then $\lim_{x \rightarrow a} f(x)/g(x) = \lim_{x \rightarrow a} f'(x)/g'(x)$*

Differentiation of inverse functions

- Implicit differentiation: To find dy/dx , first differentiate each term with respect to x treating y terms as functions of x using chain rule with dy/dx then rearrange for dy/dx
 - If $x^2 + y^2 = 1$, $2x + 2y \cdot dy/dx = 0$ and $2x \cdot dx/dy + 2y = 0$
- Let $y = f(x)$, then $[f^{-1}(y)]' = [f'(x)]^{-1}$ then rewrite in terms of y (e.g. by using below instead)

Proof: By implicit differentiation, $1 = f'(x) \cdot dx/dy$. So, $dx/dy = 1/f'(x)$

- Corollary: Let $y = f(x)$, then $[f^{-1}(y)]' \equiv [f'(f^{-1}(y))]^{-1}$ — **useful for differentiating inverse trig functions**

Alternate Proof: We have that $f(f^{-1}(y)) \equiv y$

Differentiating both sides using chain rule: $[(f^{-1}(y))]' \cdot [f'(f^{-1}(y))] = 1$

- To find $d/dx(\arcsin(x))$:

Let $f(x) = \sin(x)$

Then, $f^{-1}(x) = \arcsin(x)$ and $f'(x) = \cos(x)$

Thus, $d/dx(\arcsin(x)) = 1/\cos(\arcsin(x))$

Let $a = \arcsin(x)$

Then, $\sin^2(a) + \cos^2(a) \equiv 1$ and $x = \sin(a)$

Thus, $d/dx(\arcsin(x)) = 1/\sqrt{1 - x^2}$

Taylor series

Taylor series

- **Taylor series of $f(x)$ at $x=a$: $f(x) = \sum f^{(i)}(a)/n! \cdot (x-a)^i$ (if f can be differentiated any number of times)**
- Taylor polynomial: $f(x) = \sum f^{(i)}(a)/n! \cdot (x-a)^i$ from $i=0$ to $i=n$ (+ $R_n(x)$ where $R_n(x)$ is the difference between $f(x)$ and the Taylor polynomial)
 - **Taylor's theorem: If $f^{(n+1)}$ exists over $(a; x)$ and $f^{(n)}$ is continuous over $[a; x]$, then $\exists c \in (a; x): R_n = f^{(n+1)}(c)/(n+1)! \cdot (x-a)^{n+1}$ and $\lim_{x \rightarrow a} R_n(x) = 0$**
 - **For x in the radius of convergence, $\lim_{n \rightarrow \infty} R_n(x) = 0$**
- For x in the radius of convergence, $a_n = f^{(i)}(a)/n! \cdot (x-a)^i$ passes ratio test for $n \rightarrow \infty$
- Taylor series are power series!
- *Let $g(x) = \sum f^{(i)}(a)/n! \cdot (x-a)^i$ from $i=0$ to $i=n$. Then, $g(a)=f(a)$ and for all $N \leq n$ $g^{(N)}(a)=f^{(N)}(a)$*
- Maclaurin series = Taylor series at $a=0$

Integration

Integration

- Let $\{x_0, \dots, x_n\}$ be a partition of $[a; b]$ (a subset of $[a; b]$ s.t. $x_0=a$, $x_n=b$, and the elements are strictly increasing)
 - Let $m_r = \min \{f(x): x \in [x_{r-1}, x_r]\}$ and $M_r = \max \{f(x): x \in [x_{r-1}, x_r]\}$
 - Let $L = \sum (x_r - x_{r-1})m_r$ from $r=1$ to $r=n$ and $U = \sum (x_r - x_{r-1})M_r$ from $r=1$ to $r=n$
 - Then, L is the area formed by tallest rectangles that lie entirely below f and U is the area formed by shortest rectangles that lie entirely above f
 - f is integrable over $[a; b]$ (and A is the definite integral of f between a and b) iff there exists exactly one A : $L \leq A \leq U$ for every partition of $[a; b]$
- **The following are each necessary and sufficient for f is integrable over $[a;b]$:**
 - **f is continuous over $[a;b]$**
 - **f is increasing over $[a;b]$**
 - **f is decreasing over $[a;b]$**
- First fundamental theorem of calculus: Let $f: [a; b] \mapsto \mathbb{R}$ and $F(x) = \int f(t) dt$ from $t=a$ to $t=x$. **If f is continuous at $c \in [a;b]$, then F is differentiable at c and moreover $F'(c)=f(c)$**
- Second fundamental theorem of calculus: Let $f: [a; b] \mapsto \mathbb{R}$ and F be a function s.t. $\forall x \in [a;b]$ $F'(x)=f(x)$. Then, $\int f(x) dx$ from b to $a = F(b) - F(a)$

(First order ordinary) Differential equations

Terminology

- Ordinary DE (ODE) = DE where all derivatives are with respect to the same variable
- Order of a DE = order of highest order derivative it contains
- In this module we only deal with first order ODEs
 - Separable iff can be written in the form $dy/dx = f(x)g(y)$
 - Then $\int 1/g(y) dy = \int f(x) dx$
 - **Homogenous iff can be written in form $dy/dx=f(y/x)$**
 - Making the **substitution $v=y/x$** gives **$\log(x) = \int 1/(f(v)-v) dv$** — reverse substitution once evaluated
Proof: We have that $y=vx$
Implicit differentiation with respect to x : $dy/dx = v+x \cdot dv/dx$
 $dv/dx = (dy/dx - v)/x$
 $dv/dx = (f(v) - v)(1/x)$ — separable
 - **Linear iff can be written in the form $dy/dx+P(x)y=Q(x)$**

Solving linear (the integrating factor)

- *If $Q=0$ for all x , then is separable ($f(x)=P(x)$ and $g(y)=1/y$) and $y=e^{\int -P(x) dx}$*
- Let $I(x) = e^{\int P(x) dx}$, then $I'(x)=P(x)I(x)$
- Multiplying both sides of DE by $I(x)$ gives $I(x)y' + I'(x)y = I(x)Q(x)$
 - By the product rule, $d/dx(I(x)y) = I'(x)y + I(x)y'$. So, $d/dx(I(x)y) = I(x)Q(x)$
 - So, $I(x)y = \int I(x)Q(x) dx$
 - **Don't need +C in $I(x)$ but do need it here** and don't forget to divide it by $I(x)$
 - Thus, $y = 1/I(x)[\int I(x)Q(x) dx]$