

CS 577 - Divide and Conquer

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DIVIDE AND CONQUER

DIVIDE AND CONQUER (DC)

Overview

- Split problem into smaller sub-problems.
- Solve (usually recurse on) the smaller sub-problems.
- Use the output from the smaller sub-problems to build the solution.

Tendencies of DC

- Naturally recursive solutions
- Solving complexities often involve recurrences.
- Often used to improve efficiency of efficient solutions, e.g. $O(n^2) \rightarrow O(n \log n)$.
- Used in conjunction with other techniques.

SEARCHING

Linear Search

- Brute force approach: check every item in order.
- Time complexity: $O(n)$

Divide and Conquer Approach

- Binary Search
- Complexity: $O(\log n)$

MERGESORT

SORTING

Ordering some (multi)set of n items.

Brute Force

- Test all possible orderings.
- $O(n \cdot n!)$

Simple Sorts

- Insertion Sort, Selection Sort, Bubble Sort
- $O(n^2)$

Efficient Sorts

- Divide & Conquer: Quick Sort ($O(n^2)$), Merge Sort ($O(n \log n)$)

SORTING

Ordering some (multi)set of n items.

Simple Sorts

- Insertion Sort, Selection Sort, Bubble Sort
- $O(n^2)$

Efficient Sorts

- Divide & Conquer: Quick Sort ($O(n^2)$), Merge Sort ($O(n \log n)$)

Trick Sorts

- Radix Sort ($O(n \lceil \log k \rceil)$), Counting Sort ($O(n + k)$)
- k is the maximum key size.
- TopHat 5: What value of k would make both sorts have time complexity no better than Merge Sort? $\Omega(n \log n)$

MERGESORT

Algorithm: MERGESORT

Input : A list A of n comparable items.

Output: A sorted list A .

if $|A| = 1$ **then return** A

$A_1 := \text{MERGESORT}(\text{Front-half of } A)$

$A_2 := \text{MERGESORT}(\text{Back-half of } A)$

return $\text{MERGE}(A_1, A_2)$

Algorithm: MERGE

Input : Two lists of comparable items: A and B .

Output: A merged list.

Initialize S to an empty list.

while *either A or B is not empty* **do**

 | Pop and append $\min\{\text{front of } A, \text{front of } B\}$ to S .

end

return S

TopHat 6: What is the complexity of MERGE? $O(n)$

MERGESORT

Algorithm: MERGESORT

Input : A list A of n comparable items.

Output: A sorted list A .

if $|A| = 1$ **then return** A

$A_1 := \text{MERGESORT}(\text{Front-half of } A)$

$A_2 := \text{MERGESORT}(\text{Back-half of } A)$

return $\text{MERGE}(A_1, A_2)$

Program Correctness:

- ① **Soundness:** List A is sorted after call to MERGESORT.

Proof: By strong induction on list length:

Base case: $k = 1$: List is sorted.

Inductive step: By ind hyp, A_1 and A_2 are sorted, and, then, by definition, MERGE will produce a sorted list.

MERGESORT

Algorithm: MERGESORT

Input : A list A of n comparable items.

Output: A sorted list A .

if $|A| = 1$ **then return** A

$A_1 := \text{MERGESORT}(\text{Front-half of } A)$

$A_2 := \text{MERGESORT}(\text{Back-half of } A)$

return $\text{MERGE}(A_1, A_2)$

Program Correctness:

- 1 Soundness: List A is sorted after call to MERGESORT.
- 2 Complete: Handles lists of any size, and each recursion makes progress towards base case by splitting the list in half.

MERGESORT

Algorithm: MERGESORT

Input : A list A of n comparable items.

Output: A sorted list A .

if $|A| = 1$ **then return** A

$A_1 := \text{MERGESORT}(\text{Front-half of } A)$

$A_2 := \text{MERGESORT}(\text{Back-half of } A)$

return $\text{MERGE}(A_1, A_2)$

Run time Considerations:

- Cost to MERGE: $O(n)$.
- Recurrences: 2 calls to MERGESORT with lists half the size.

MERGESORT RECURRENCE

$$T(n) \leq 2 \cdot T\left(\frac{n}{2}\right) + cn; T(1) \leq c$$

Notes

- More precise: $T(n) \leq T\left(\lfloor \frac{n}{2} \rfloor\right) + T\left(\lceil \frac{n}{2} \rceil\right) + cn$
- Usually, we can asymptotically ignore floor and ceilings.
- Essentially, we are assuming n is a power of 2.
- Alternate form: $T(n) \leq 2 \cdot T\left(\frac{n}{2}\right) + O(n); T(1) \leq O(1)$

Methods

- Unwind / Recurrence Tree
- Guess
- Master Theorem
- Nuclear Bomb Theorem / Master Master Theorem

UNWIND MERGESORT RECURRENCE

$$T(n) \leq 2T\left(\frac{n}{2}\right) + cn$$

$$\leq 2\left(2T\left(\frac{n}{4}\right) + c\frac{n}{2}\right) + cn$$

$$\leq 2\left(2\left(2T\left(\frac{n}{2^3}\right) + c\frac{n}{2^2}\right) + c\frac{n}{2}\right) + cn$$

$$\vdots$$

$$\leq 2^k T\left(\frac{n}{2^k}\right) + kcn$$

$$= nT(1) + cn \log(n)$$

$$= cn + cn \log n$$

$$= O(n \log(n))$$

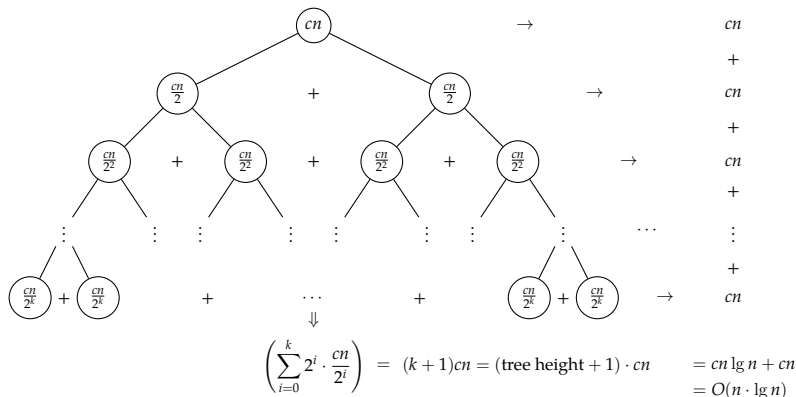
$$1 = \frac{n}{2^k}$$

$$\iff 2^k = n$$

$$\iff k = \log_2(n)$$

RECURSION TREE METHOD

$$T(n) \leq 2 \cdot T\left(\frac{n}{2}\right) + cn; T(1) \leq c$$



¹Based on: <http://www.texample.net/tikz/examples/merge-sort-recursion-tree/>

GUESS METHOD

$$T(n) \leq 2 \cdot T\left(\frac{n}{2}\right) + cn; T(1) \leq c$$

Procedure

- ➊ Guess: Seems like $O(n \log n)$ -ish.
- ➋ Prove by induction! Not valid without proof!

PROVE RECURRENCE BY STRONG INDUCTION

$$T(n) \leq 2 \cdot T\left(\frac{n}{2}\right) + cn \leq cn \lg n + cn; T(1) \leq c$$

Base Case: $n = 2$.

$$\begin{aligned} T(2) &= 2 \cdot T(1) + 2c \leq 4c \\ &= c \cdot 2 \lg 2 + 2c \end{aligned}$$

Inductive step:

$$\begin{aligned} T(k) &= 2 \cdot T(k/2) + ck \\ &\leq 2 \left(\frac{ck}{2} \lg \frac{k}{2} + \frac{ck}{2} \right) + ck \\ &= ck \lg(k/2) + 2ck \\ &= ck \lg k - ck + 2ck \\ &= ck \lg k + ck \end{aligned}$$

$\therefore O(n \log n)$

GENERALIZED RECURRENCE

$$T(n) \leq q \cdot T\left(\frac{n}{2}\right) + cn; T(1) \leq c$$

Case $q > 2$

$$O(n^{\lg q})$$

Case $q = 2$

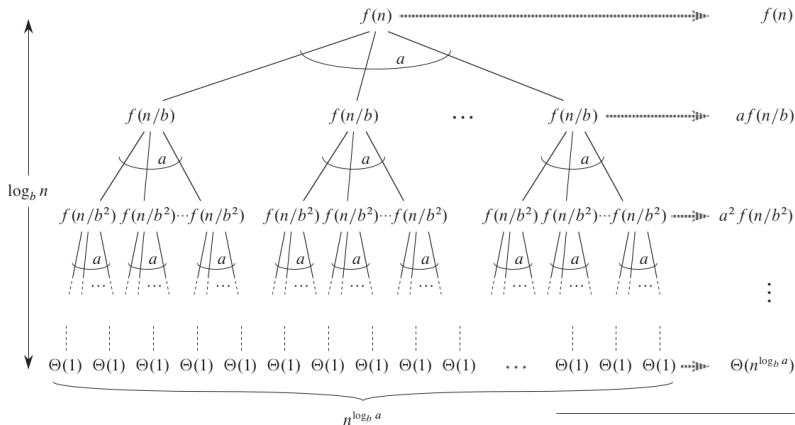
$$O(n \log n)$$

Case $q = 1$

$$O(n)$$

MASTER THEOREM

COOKBOOK RECURRENCE SOLVING



$$\text{Total: } \Theta(n^{\log_b a}) + \sum_{j=0}^{\log_b n - 1} a^j f(n/b^j)$$

MASTER THEOREM

COOKBOOK RECURRENCE SOLVING

Theorem 1

Let $a \geq 1$ and $b > 1$ be constants, let $f(n)$ be a function, and let $T(n)$ be defined on the non-negative integers by the recurrence

$$T(n) = aT(n/b) + f(n) ,$$

where we interpret n/b to mean either $\lfloor n/b \rfloor$ or $\lceil n/b \rceil$. Then $T(n)$ has the following asymptotic bounds:

- ❶ If $f(n) = O(n^{\log_b a - \varepsilon})$ for some constant $\varepsilon > 0$, then $T(n) = \Theta(n^{\log_b a})$.
- ❷ If $f(n) = \Theta(n^{\log_b a})$, then $T(n) = \Theta(n^{\log_b a} \log n)$.
- ❸ If $\Omega(n^{\log_b a + \varepsilon})$ for some constant $\varepsilon > 0$, and if $a \cdot f(n/b) \leq c \cdot f(n)$ for some constant $c < 1$ and all sufficiently large n , then $T(n) = \Theta(f(n))$.

NUCLEAR BOMB / MASTER MASTER THEOREM

AKRA AND BAZZI, 1998

Theorem 2

Given a recurrence of the form:

$$T(n) = \sum_{i=1}^k a_i T(n/b_i) + f(n) ,$$

where k is a constant, $a_i > 0$ and $b_i > 1$ are constants for all i , and $f(n) = \Omega(n^c)$ and $f(n) = O(n^d)$ for some constants $0 < c \leq d$.

Then,

$$T(n) = \Theta \left(n^{\rho} \left(1 + \int_1^n \frac{f(u)}{u^{\rho+1}} du \right) \right) ,$$

where ρ is the unique real solution to the equation

$$\sum_{i=1}^k \frac{a_i}{b_i^{\rho}} = 1 .$$

INVERSION COUNT

COUNTING INVERSIONS

Inversion

Given a list A of comparable items. An inversion is a pair of items (a_i, a_j) such that $a_i > a_j$ and $i < j$, where i and j are the index of the items in A .

Inversion Count

Count the number of inversions in a list A , containing n comparable items.

Exercise – Teams of 2 or 3

- Solve the problem in $\Theta(n^2)$.
- Solve the problem in $O(n \log n)$.
- Prove correctness and complexity.

PART 1: GIVE A $\Theta(n^2)$ SOLUTION.

Algorithm: CHECKALLPAIRS

Input : A list A of n comparable items.

Output: Number of inversions in A .

Let $c := 0$

```
for  $i := 1$  to  $\text{len}(A) - 1$  do
    for  $j := i$  to  $\text{len}(A)$  do
        if  $A[i] > A[j]$  then
             $c := c + 1$ 
        end
    end
end
return  $c$ 
```

Analysis

- Correct: Checks all pairs and counts the inversions.
- Complexity: For each i , check $n - i$ pairs. Overall:

$$\sum_{i=1}^{n-1} i = \frac{n(n-1)}{2} = \Theta(n^2).$$

PART 2: GIVE AN $O(n \log n)$ SOLUTION.

Algorithm: COUNTSORT

Input : A list A of n comparable items.

Output: A sorted array and the number of inversions.

if $|A| = 1$ **then return** $(A, 0)$

$(A_1, c_1) := \text{COUNTSORT}(\text{Front-half of } A)$

$(A_2, c_2) := \text{COUNTSORT}(\text{Back-half of } A)$

$(A, c) := \text{MERGECOUNT}(A_1, A_2)$

return $(A, c + c_1 + c_2)$

PART 2: GIVE AN $O(n \log n)$ SOLUTION.

Algorithm: MERGECOUNT

Input : Two lists of comparable items: A and B .

Output: A merged list and the count of inversions.

Initialize S to an empty list and $c := 0$.

while *either A or B is not empty* **do**

 Pop and append $\min\{\text{front of } A, \text{front of } B\}$ to S .

if *Appended item is from B* **then**

$c := c + |A|$.

end

end

return (S, c)

Analysis

- Correctness: Need to show that the inversions are counted.
- Complexity: Same recurrence as MERGESORT.

LINEAR TIME SELECTION

LINEAR TIME SELECTION

Problem

Find the k th value in an unsorted array A of n numbers if A were sorted.

Algorithm: QUICKSELECT

Input : A array $A[1..n]$ and an int k .

Output: The k th element of A if A were sorted.

if $n = 1$ **then return** $A[1]$

Choose a pivot $A[p]$

$r := \text{PARTITION}(A[1..n], p)$

if $k < r$ **then**

return $\text{QUICKSELECT}(A[1..r - 1], k)$

else if $k > r$ **then**

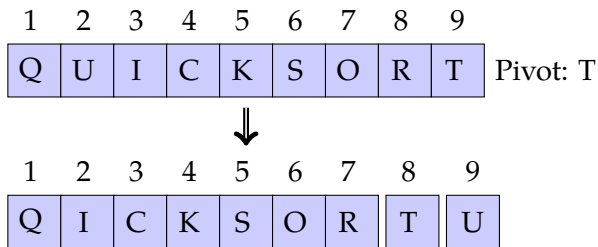
return $\text{QUICKSELECT}(A[r + 1..n], k - r)$

else

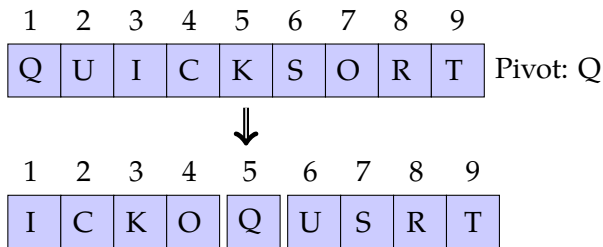
return $A[r]$

end

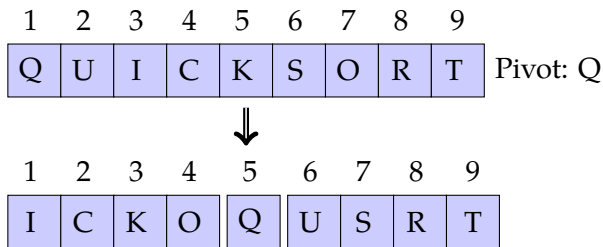
PARTITION AROUND A PIVOT



PARTITION AROUND A PIVOT



PARTITION AROUND A PIVOT



How much work is done to partition around a pivot? $O(n)$

QUICKSELECT RECURRENCE

$$T(n) \leq \max_{1 \leq r \leq n} \max\{T(r-1), T(n-r)\} + cn$$

Algorithm: QUICKSELECT

Input : A array $A[1..n]$ and an int k .

Output: The k th element of A .

if $n = 1$ **then return** $A[1]$

Choose a pivot $A[p]$

$r := \text{PARTITION}(A[1..n], p)$

if $k < r$ **then**

return $\text{QUICKSELECT}(A[1..r-1], k)$

else if $k > r$ **then**

return $\text{QUICKSELECT}(A[r+1..n], k-r)$

else

return $A[r]$

end

QUICKSELECT RECURRENCE

$$\begin{aligned} T(n) &\leq \max_{1 \leq r \leq n} \max\{T(r-1), T(n-r)\} + cn \\ &\leq \max_{1 \leq \ell \leq n-1} T(\ell) + cn \\ &\leq T(n-1) + cn \\ &\in O(n^2) \end{aligned}$$

MEDIAN OF MEDIAN

Algorithm: MOMSELECT

Input : A array $A[1..n]$ and an int k .

Output: The k th element of A .

if n is small **then** Solve by brute force.

$m := \lceil n/5 \rceil$

for $i := 1$ to m **do**

$M[i] :=$ brute force find median of $A[5i - 4..5i]$

end

$mom := \text{MOMSELECT}(M[1..m], \lfloor m/2 \rfloor)$

$r := \text{PARTITION}(A[1..n], mom)$

if $k < r$ **then**

return $\text{MOMSELECT}(A[1..r - 1], k)$

else if $k > r$ **then**

return $\text{MOMSELECT}(A[r + 1..n], k - r)$

else

return $A[r]$

end

MOMSELECT ANALYSIS

MomSelect Pivot

- greater and less than $> \lfloor \lceil n/5 \rceil / 2 \rfloor - 1 \approx n/10$ medians.
- Therefore, MomSelect Pivot is greater and less than $3n/10$ items.
- So, worst-case partition size is $7n/10$.

Recurrence:

$$T(n) \leq T(n/5) + T(7n/10) + cn \in O(n)$$

INTEGER MULTIPLICATION

INTEGER MULTIPLICATION

Partial Product Method:

$$\begin{array}{r}
 12 \\
 \times 13 \\
 \hline
 36 \\
 12 \\
 \hline
 156
 \end{array}
 \qquad
 \begin{array}{r}
 1100 \\
 \times 1101 \\
 \hline
 1100 \\
 0000 \\
 1100 \\
 1100 \\
 \hline
 10011100
 \end{array}$$

Problem

Multiply two binary numbers x and y , counting every bitwise operation.

TopHat 8: What is the complexity of the partial product method? $O(n^2)$.

DIVIDE & CONQUER v1

High and low bits

Consider $x = x_1 \cdot 2^{n/2} + x_0$ and $y = y_1 \cdot 2^{n/2} + y_0$.

$$\begin{aligned} xy &= (x_1 \cdot 2^{n/2} + x_0)(y_1 \cdot 2^{n/2} + y_0) \\ &= x_1 y_1 \cdot 2^n + (x_1 y_0 + x_0 y_1) \cdot 2^{n/2} + x_0 y_0 \end{aligned}$$

- How many recursive calls? 4.
- Cost per call? $O(n)$
- What is the size of the recursive calls? $n/2$.
- What is the recurrence?

$$T(n) \leq 4T(n/2) + cn = O\left(n^{\lg 4}\right) = O\left(n^2\right)$$

DIVIDE & CONQUER v2

High and low bits

Consider $x = x_1 \cdot 2^{n/2} + x_0$ and $y = y_1 \cdot 2^{n/2} + y_0$.

$$\begin{aligned} xy &= (x_1 \cdot 2^{n/2} + x_0)(y_1 \cdot 2^{n/2} + y_0) \\ &= x_1y_1 \cdot 2^n + (x_1y_0 + x_0y_1) \cdot 2^{n/2} + x_0y_0 \end{aligned}$$

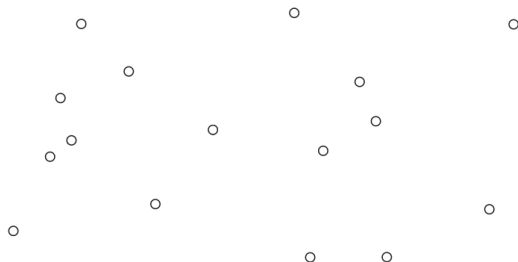
Hint: $(x_1 + x_0)(y_1 + y_0) = x_1y_1 + x_1y_0 + x_0y_1 + x_0y_0$.

Exercise: Design an algorithm with 3 Recursive Calls

- Recursions:
 - $p := \text{intMult}(x_1 + x_0, y_1 + y_0)$
 - $x_1y_1 := \text{intMult}(x_1, y_1)$
 - $x_0y_0 := \text{intMult}(x_0, y_0)$
- Combine: Return $x_1y_1 \cdot 2^n + (p - x_1y_1 - x_0y_0) \cdot 2^{n/2} + x_0y_0$
- Recurrence: $T(n) \leq 3T(n/2) + O(n) = O(n^{\lg 3}) = O(n^{1.59})$

CLOSEST PAIRS

FINDING THE CLOSEST PAIR OF POINTS



Problem

Given a set of n points, $\mathcal{P} = \{p_1, p_2, \dots, p_n\}$, in the plane. Find the closest pair. That is, solve $\arg \min_{(p_i, p_j) \in \mathcal{P}} \{d(p_i, p_j)\}$, where $d(\cdot, \cdot)$ is the Euclidean distance.

What is the $O(n^2)$ solution?

1-D VERSION

1-d Closest Pair

The points are on the line.

$O(n \log n)$ for 1-d Closest Pair

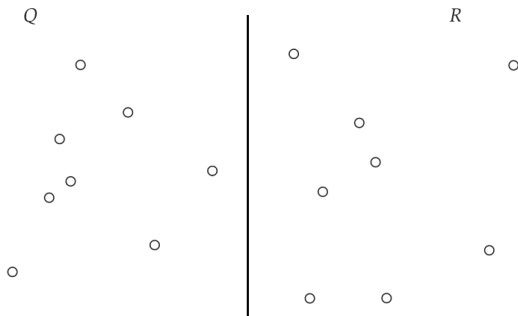
- Sort the points ($O(n \log n)$).
- Walk through sorted points and find minimum pair ($O(n)$).

2-D CLOSEST PAIR

DIVIDE AND CONQUER

- ➊ Divide: Split point set (in half?).
- ➋ Conquer: Find closest pair in each partition.
- ➌ Combine: Merge the solutions.

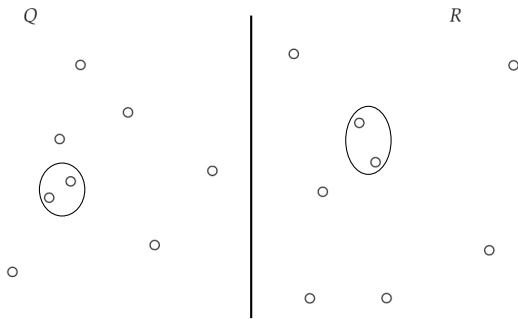
1. DIVIDE: SPLIT THE POINTS



Definitions

- \mathcal{P}_x : Points sorted by x -coordinate.
- \mathcal{P}_y : Points sorted by y -coordinate.
- Q (resp. R) is left (resp. right) half of \mathcal{P}_x .

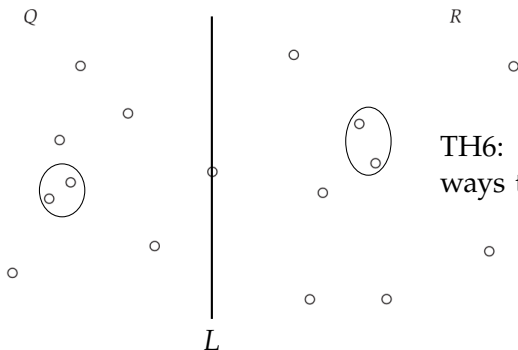
2. CONQUER: FIND THE MIN IN Q AND R



Key Observations

- From \mathcal{P}_x and \mathcal{P}_y : We can create Q_x, Q_y, R_x, R_y without resorting.
- Running time for this: $O(n)$.
- Let (q_0^*, q_1^*) and (r_0^*, r_1^*) be closest pairs in Q and R .

3. COMBINE THE SOLUTIONS.

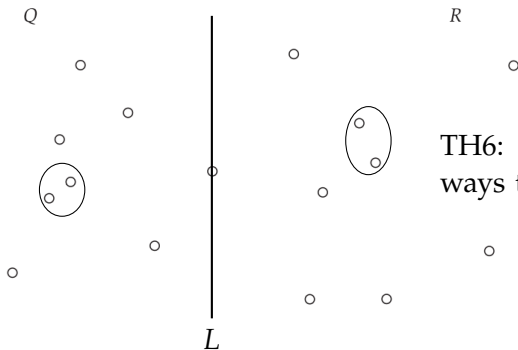


TH6: Are one of these always the minimum of \mathcal{P} ?

Claim 1

Let $\delta := \min\{d(q_0^*, q_1^*), d(r_0^*, r_1^*)\}$. If there exists a $q \in Q$ and an $r \in R$ for which $d(q, r) < \delta$, then each of q and r are within δ of L .

3. COMBINE THE SOLUTIONS.



TH6: Are one of these always the minimum of \mathcal{P} ?

Lemma 3

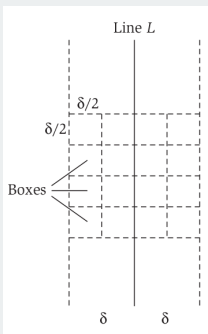
Let S be the set of points within δ of L . If there exists a $s, s' \in S$ and $d(s, s') < \delta$, then s and s' are within 15 positions of each other in S_y .

3. COMBINE THE SOLUTIONS.

Lemma 3

Let S be the set of points within δ of L . If there exists a $s, s' \in S$ and $d(s, s') < \delta$, then s and s' are within 15 positions of each other in S_y .

Proof.



- Partition δ -space around L into $\delta/2$ squares.
- At most 1 point per square else contradicts definition of δ .
- By way of contradiction, say $d(s, s') < \delta$ and s and s' separated by 16 positions.
- By counting argument, s and s' are separated by 3 rows which is at least $3\delta/2$. □

3. COMBINE THE SOLUTIONS.

Lemma 3

Let S be the set of points within δ of L . If there exists a $s, s' \in S$ and $d(s, s') < \delta$, then s and s' are within 15 positions of each other in S_y .

Completing the Algorithm

- Find the min pair (s, s') in S .
 - For each $p \in S$, check the distance to each of next 15 points in S_y .
- If $d(s, s') < \delta$, return (s, s')
- else return min of (q_0^*, q_1^*) and (r_0^*, r_1^*) .

COMPLETING THE ANALYSIS

Correctness of the Algorithm

- By induction on the number of points.
- Use the definition of the algorithm and the claims establish in Step 3.

Runtime of the Algorithm

- Sorting by x and by y ($O(n \log n)$).
- How many recursive calls? 2.
- What is the size of the recursive calls? $n/2$.
- Work per call: check points in S .
 - $15 \cdot |S| = O(n)$.
- What is the recurrence?

$$T(n) \leq 2T(n/2) + cn = O(n \log n) .$$

MAX SUBARRAY

MAX SUBARRAY

Problem

Given an array A of integers, find the (non-empty) contiguous subarray of A of maximum sum.

Exercise – Teams of 3 or so

- Solve the problem in $\Theta(n^2)$.
- Solve the problem in $O(n \log n)$.
- Prove correctness and complexity.

PART 1: GIVE A $\Theta(n^2)$ SOLUTION.

Algorithm: CHECKALLSUBARRAYS

Input : Array A of n ints.**Output:** Max subarray in A .Let M be an empty array

```
for  $i := 1$  to  $\text{len}(A)$  do
  for  $j := i$  to  $\text{len}(A)$  do
    if  $\text{sum}(A[i..j]) > \text{sum}(M)$ 
      |  $M := A[i..j]$ 
    end
  end
end
return  $M$ 
```

Analysis

- Correct: Checks all possible contiguous subarrays.
- Complexity:
 - Re-calculating the sum will make it $O(n^3)$. Key is to calculate the sum as you iterate.
 - For each i , check $n - i + 1$ ends. Overall:

$$\sum_{i=1}^n i = \frac{n(n+1)}{2} = \Theta(n^2)$$

PART 2: GIVE AN $O(n \log n)$ SOLUTION.

Algorithm: MAXSUBARRAY

Input : Array A of n ints.

Output: Max subarray in A .

if $|A| = 1$ **then return** $A[1]$

$A_1 := \text{MAXSUBARRAY}(\text{Front-half of } A)$

$A_2 := \text{MAXSUBARRAY}(\text{Back-half of } A)$

$M := \text{MIDMAXSUBARRAY}(A)$

return *Array with max sum of $\{A_1, A_2, M\}$*

Algorithm: MIDMAXSUBARRAY

Input : Array A of n ints.

Output: Max subarray that crosses midpoint A .

$m := \text{mid-point of } A$

$L := \text{max subarray in } A[i, m-1] \text{ for } i = m-1 \rightarrow 1$

$R := \text{max subarray in } A[m, j] \text{ for } j = m \rightarrow n$

return $L \cup R$ // subarray formed by combining L and R .

PART 2: GIVE AN $O(n \log n)$ SOLUTION.

Algorithm: MAXSUBARRAY

Input : Array A of n ints.

Output: Max subarray in A .

if $|A| = 1$ **then return** $A[1]$

$A_1 := \text{MAXSUBARRAY}(\text{Front-half of } A)$

$A_2 := \text{MAXSUBARRAY}(\text{Back-half of } A)$

$M := \text{MIDMAXSUBARRAY}(A)$

return Array with max sum of $\{A_1, A_2, M\}$

Analysis

- Correctness: By induction, A_1 and A_2 are max for subarray and M is max mid-crossing array.
- Complexity: Same recurrence as MERGESORT.

MATRIX MULTIPLICATION

MATRIX MULTIPLICATION

Problem

Multiple two $n \times n$ matrices, A and B . Let $C = AB$.

$$\begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix} \begin{bmatrix} 4 & 3 \\ 2 & 1 \end{bmatrix} = \begin{bmatrix} 1 \cdot 4 + 2 \cdot 2 & 1 \cdot 3 + 2 \cdot 1 \\ 3 \cdot 4 + 4 \cdot 2 & 3 \cdot 3 + 4 \cdot 1 \end{bmatrix} = \begin{bmatrix} 8 & 5 \\ 20 & 13 \end{bmatrix}$$

Algorithm: Naïve Method

```
for  $i \leftarrow 1$  to  $n$  do
  for  $j \leftarrow 1$  to  $n$  do
    for  $k \leftarrow 1$  to  $n$  do
       $C[i][j] += A[i][k] \cdot B[k][j]$ 
    end
  end
end
```

TopHat 12: What is the complexity of the Naïve Method? $O(n^3)$.

DIVIDE & CONQUER V1

$$\left[\begin{array}{c|c} a & b \\ \hline c & d \end{array} \right] \left[\begin{array}{c|c} e & f \\ \hline g & h \end{array} \right] = \left[\begin{array}{c|c} ae + bg & af + bh \\ \hline ce + dg & cf + dh \end{array} \right]$$

- How many recursive calls? 8.
- Cost per call? $O(n^2)$ time per addition
- What is the size of the recursive calls? $n/2$.
- What is the recurrence?

$$T(n) \leq 8T(n/2) + cn^2 = O\left(n^{\lg 8}\right) = O\left(n^3\right)$$

DIVIDE & CONQUER v2

$$\left[\begin{array}{c|c} a & b \\ \hline c & d \end{array} \right] \left[\begin{array}{c|c} e & f \\ \hline g & h \end{array} \right] = \left[\begin{array}{c|c} \frac{p_5 + p_4 - p_2 + p_6}{p_3 + p_4} & \frac{p_1 + p_2}{p_1 + p_5 - p_3 - p_7} \end{array} \right]$$

Strassen's Method (1969)

- $p_1 := a(f - h)$
- $p_2 := (a + b)h$
- $p_3 := (c + d)e$
- $p_4 := d(g - e)$
- $p_5 := (a + d)(e + h)$
- $p_6 := (b - d)(g + h)$
- $p_7 := (a - c)(e + f)$

What is the recurrence?

$$T(n) \leq 7T(n/2) + cn^2 = O\left(n^{\lg 7}\right) = O\left(n^{2.8074}\right)$$

DIVIDE & CONQUER v2

$$\left[\begin{array}{c|c} a & b \\ \hline c & d \end{array} \right] \left[\begin{array}{c|c} e & f \\ \hline g & h \end{array} \right] = \left[\begin{array}{c|c} \frac{p_5 + p_4 - p_2 + p_6}{p_3 + p_4} & \frac{p_1 + p_2}{p_1 + p_5 - p_3 - p_7} \end{array} \right]$$

Current Champ: $O(n^{2.373})$



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