

Exercise 5.1

Of which Chomsky type are the following productions?

(1) $B \rightarrow cA$ = Type: 0, 2, 3

(2) $cA \rightarrow B$ = Type: 0

(3) $cA \rightarrow BaB$ = Type: 0

(4) $C \rightarrow aBc$ = Type: 0, 2

0: $(NT \text{ or terminal})^* \cdot NT(NT \text{ or terminal})^* \rightarrow (NT \text{ or terminal})^*$

1: $\sigma NT \sigma \rightarrow \sigma(NT \text{ or terminal}) + \sigma$

$\sigma = (NT \text{ or terminal})^*$

2: $(NT) \rightarrow (NT \text{ or terminal})^*$

3: $(NT) \rightarrow (\text{terminal to } NT)$

$(NT) \rightarrow (\text{terminal})$

$(NT) \rightarrow (\epsilon)$

Exercise 5.2 - obligatory (4 points)

The following DTD (document type definition) for XML documents is given:

```
<!DOCTYPE a [
```

```
<!ELEMENT a (b | c)>
```

```
<!ELEMENT b (c, d?)*>
```

```
<!ELEMENT c (#PCDATA)>
```

```
<!ELEMENT d (#PCDATA)>
```

```
]>
```

*(b or c)
(c, d is optional) c is optional with ε*

Which of the following XML documents are valid with respect to this DTD? Indicate all errors which are contained in the documents.

(1) Valid

```
<a>
```

```
<c>
```

```
xyz
```

```
</c>
```

```
</a>
```

<c> should contain only elementary text

(2) Invalid

```
<a>
```

```
<b>
```

```
<c> 12 </c>
```

```
<d> 34 </d>
```

```
<d> 56 </d>
```

```
</b>
```

```
</a>
```

- the second element <d> is not possible. A <c> must follow.

(3) Invalid

```
<a>
```

```
<b>
```

```
<d> 555 </d>
```

```
<c> 444 </c>
```

```
<d> 333 </d>
```

```
</b>
```

(4) Valid

```
<a>
```

```
<b>
```

```
<c> rrr </c>
```

<c>666</c>

- element b should be 2nd to element c.

- a should only contain elements a or b.

<c>sss</c>

<d>ttt</d>

<c>uuu</c>

- the child elements of <d> are in the correct position.

Exercise 5.3 - obligatory (6 points)

Let $A = (Q, \Sigma, \delta, z_0, E)$ be a deterministic finite automaton (DFA), where

$Z = \{z_0, z_1, z_2, z_3, z_4, z_5\}$

$S = \{a, b\}$

$\delta(z_0, a) = z_1$ $\delta(z_0, b) = z_2$

$\delta(z_1, a) = z_1$ $\delta(z_1, b) = z_3$

$\delta(z_2, a) = z_5$ $\delta(z_2, b) = z_2$

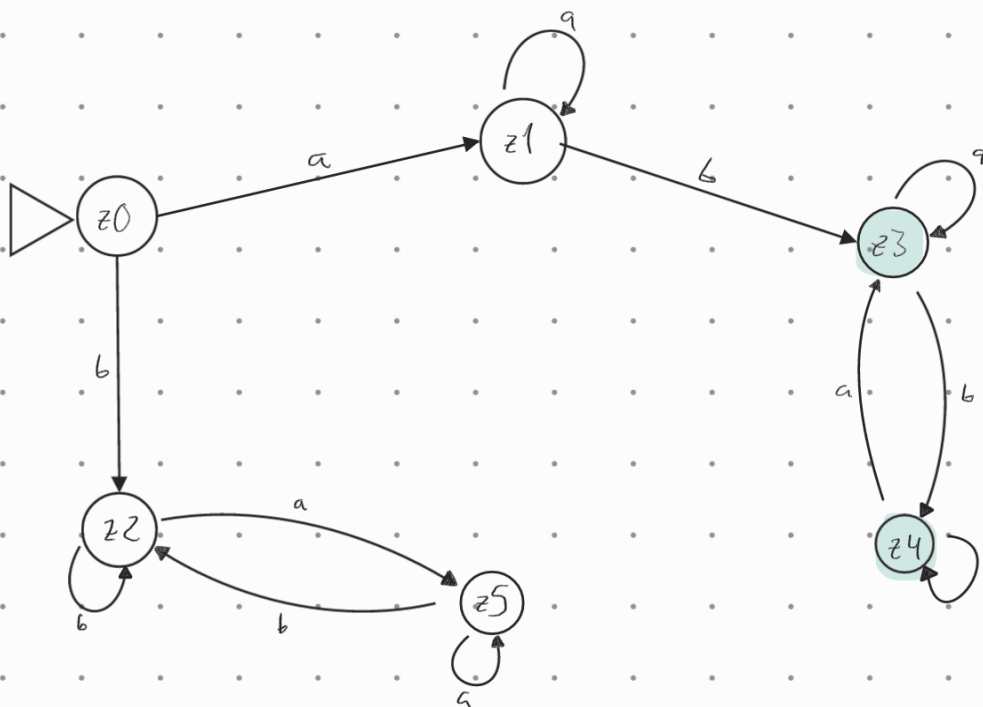
$\delta(z_3, a) = z_3$ $\delta(z_3, b) = z_4$

$\delta(z_4, a) = z_3$ $\delta(z_4, b) = z_4$

$\delta(z_5, a) = z_5$ $\delta(z_5, b) = z_2$

$E = \{z_3, z_4\}$

a) Draw A as a transition diagram.



b) Which of the following strings are accepted by A.

(1) ba = not accepted

(2) bbb = not accepted

(3) baabab = not accepted

(4) abababbbaababbba = accepted, stops at z_3

c) Which language $L(A)$ is accepted by A?

$L(A) = (a^+, b, a^*, b^+, a^+)$

$S \rightarrow AB$

$A \rightarrow aA \mid \epsilon$

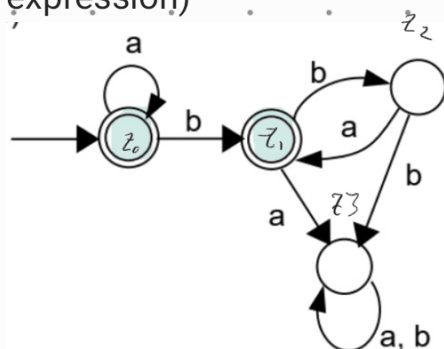
$B \rightarrow bC$

$C \rightarrow aC \mid bD \mid \epsilon$

$D \rightarrow bD \mid aC \mid \epsilon$

Exercise 5.4

What language does the following DEA accept? (can be specified as a regular expression)



$Z = \{z_0, z_1, z_2, z_3\}$

$S = \{a, b\}$

$\delta(z_0, a) = z_0$ $\delta(z_0, b) = z_1$

$\delta(z_1, a) = z_3$ $\delta(z_1, b) = z_2$

$\delta(z_2, a) = z_1$ $\delta(z_2, b) = z_3$

$\delta(z_3, a) = z_3$ $\delta(z_3, b) = z_3$

$E = \{z_0, z_1\}$

$S \rightarrow AB$

$A \rightarrow aA \mid \epsilon$

$B \rightarrow b(ba)^* \mid \epsilon$

Exercise 5.5

a) Let $S = \{0, 1\}$. Define a deterministic finite automaton (DFA) that accepts the language

$L_1 = \{w \in S^* \mid |w| \text{ is odd}\}$

$Z = \{z_0, z_1, z_2, z_3\}$

$S = \{0, 1\}$

$\delta(z_0, 1) = z_1$ $\delta(z_0, 0) = z_2$

$\delta(z_1, 0) = z_3$ $\delta(z_1, 1) = z_0$ $\delta(z_2, 0) = z_0$ $\delta(z_2, 1) = z_3$

$\delta(z_3, 0) = z_1$ $\delta(z_3, 1) = z_2$

$E = \{z_3\}$





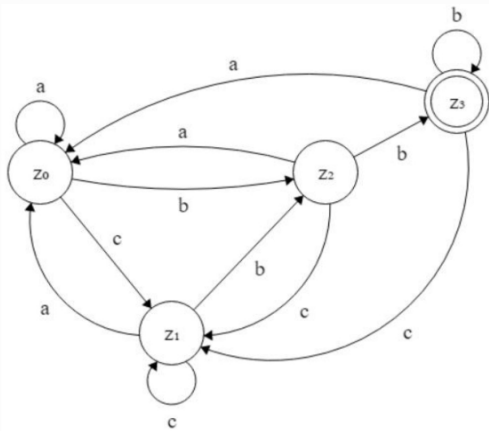
b) Let $S = \{0,1\}$. Define a DFA that accepts the language $L2 = \{11w00 \mid w \in S^*\} \cup \{00w11 \mid w \in S^*\}$.

c) Let $S = \{a, b\}$. Define a DFA accepting all strings $w \in S^*$ that start with character b and contain an odd number of a characters.

Exercise 5.6 - obligatory (6 points)

Let $S = \{a, b, c\}$.

a) Define a DFA that accepts all strings ending with bb .



$A = (Z, S, \delta, z_0, E)$

$Z = \{z_0, z_1, z_2, z_3\}$

$S = \{a, b, c\}$

$\delta : Z \times \Sigma \rightarrow Z$

$\delta(z_0, a) = z_0$

$\delta(z_0, b) = z_2$

$\delta(z_0, c) = z_1$

$\delta(z_1, a) = z_0$

$\delta(z1, b) = z2$

$\delta(z1, c) = z1$

$\delta(z2, a) = z0$

$\delta(z2, b) = z3$

$\delta(z2, c) = z1$

$\delta(z3, a) = z0$

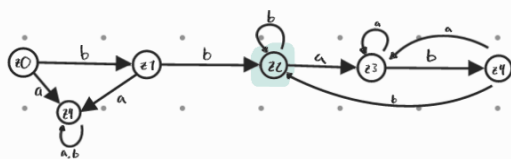
$\delta(z3, b) = z3$

$\delta(z3, c) = z1$

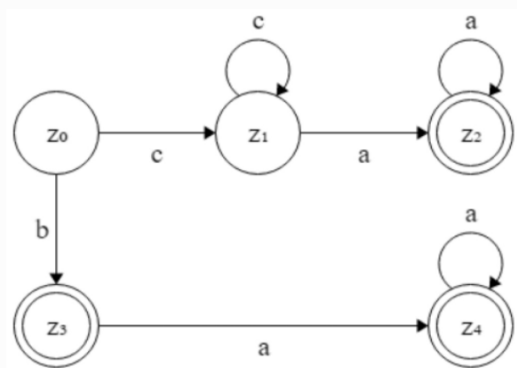
$z0$ is initial state

$E = \{z3\}$

Ending with bb with $\{a, b\}$



b) Define a DFA that accepts the language of the regular expression $(c+a|b)a^*$.



$B = (Z, S, \delta, z0, E)$

$Z = \{z0, z1, z2, z3, z4\}$

$S = \{a, b, c\}$

$\delta : Z \times \Sigma \rightarrow Z$

$\delta(z0, c) = z1$

$\delta(z1, c) = z1$

$\delta(z1, a) = z2$

$\delta(z2, a) = z2$

$\delta(z0, b) = z3$

$\delta(z3, a) = z4$

$\delta(z4, a) = z4$

$z0$ is initial state

$E = \{z2, z3, z4\}$

Exercise 5.7

Define a DFA that accepts all floating-point numbers that are build as following:

- The integer part preceding the decimal point and the fractional part after the

decimal point can consist of an arbitrary number of digits (one or more).

- If there is an exponent, it begins with 'e' or 'E', has an optional sign '+' or '-' and an arbitrary number of digits (one or more).

Here are some examples:

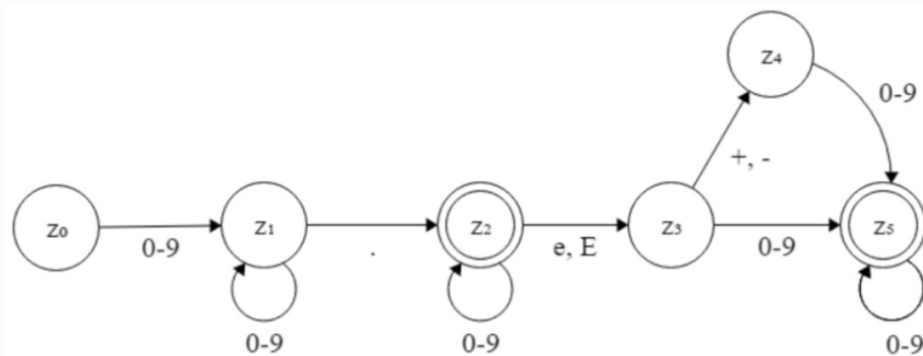
123.4500

12.345e6

0.20E+678

1004.5e-6789

You can omit error states.



$A = (Z, S, \delta, z_0, E)$

$Z = \{z_0, z_1, z_2, z_3, z_4, z_5\}$

$S = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, ., +, -, e, E\}$

$\delta : Z \times \Sigma \rightarrow Z$

$\delta(z_0, 0-9) = z_1$

$\delta(z_1, 0-9) = z_1$

$\delta(z_1, .) = z_2$

$\delta(z_2, 0-9) = z_2$

$\delta(z_2, \{e, E\}) = z_3$

$\delta(z_3, \{+, -\}) = z_4$

$\delta(z_3, 0-9) = z_5$

$\delta(z_5, 0-9) = z_5$

z_0 is initial state

$E = \{z_2, z_5\}$

