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S. Emery  
Oracle  
N. Williams  
Cryptonector  
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A Pseudo-Random Function (PRF) for the Kerberos V Generic Security  
Service Application Program Interface (GSS-API) Mechanism

Abstract

This document defines the Pseudo-Random Function (PRF) for the Kerberos V mechanism for the Generic Security Service Application Program Interface (GSS-API), based on the PRF defined for the Kerberos V cryptographic framework, for keying application protocols given an established Kerberos V GSS-API security context.

This document obsoletes [RFC 4402](#) and reclassifies that document as Historic. [RFC 4402](#) starts the PRF+ counter at 1; however, a number of implementations start the counter at 0. As a result, the original specification would not be interoperable with existing implementations.

Status of This Memo

This is an Internet Standards Track document.

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## 1. Introduction

This document specifies the Kerberos V GSS-API mechanism's [[RFC4121](#)] pseudo-random function corresponding to [[RFC4401](#)]. The function is a "PRF+" style construction. For more information, see [[RFC4401](#)], [[RFC2743](#)], [[RFC2744](#)], and [[RFC4121](#)].

This document obsoletes [RFC 4402](#) and reclassifies that document as Historic. [RFC 4402](#) starts the PRF+ counter at 1; however, a number of implementations start the counter at 0. As a result, the original specification would not be interoperable with existing implementations.

## 2. Conventions Used in This Document

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [[RFC2119](#)].

### 3. Kerberos V GSS Mechanism PRF

The GSS-API PRF [RFC4401] function for the Kerberos V mechanism [RFC4121] shall be the output of a PRF+ function based on the encryption type's PRF function keyed with the negotiated session key of the security context corresponding to the 'prf\_key' input parameter of GSS\_Pseudo\_random().

This PRF+ MUST be keyed with the key indicated by the 'prf\_key' input parameter as follows:

- o GSS\_C\_PRF\_KEY\_FULL -- use the sub-session key asserted by the acceptor (if any exists), or the sub-session asserted by the initiator (if any exists), or the Ticket's session key.
- o GSS\_C\_PRF\_KEY\_PARTIAL -- use the sub-session key asserted by the initiator (if any exists) or the Ticket's session key.

The PRF+ function is a simple counter-based extension of the Kerberos V pseudo-random function [RFC3961] for the encryption type of the security context's keys:

$$\text{PRF+}(K, L, S) = \text{truncate}(L, T0 \parallel T1 \parallel \dots \parallel Tn)$$
$$Tn = \text{pseudo-random}(K, n \parallel S)$$

where K is the key indicated by the 'prf\_key' parameter, '||' is the concatenation operator, 'n' is encoded as a network byte order 32-bit unsigned binary number, truncate(L, S) truncates the input octet string S to length L, and pseudo-random() is the Kerberos V pseudo-random function [RFC3961].

The maximum output size of the Kerberos V mechanism's GSS-API PRF then is, necessarily, 2<sup>32</sup> times the output size of the pseudo-random() function for the encryption type of the given key.

When the input size is longer than 2<sup>14</sup> octets as per [RFC4401] and exceeds an implementation's resources, then the mechanism MUST return GSS\_S\_FAILURE and GSS\_KRB5\_S\_KG\_INPUT\_TOO\_LONG as the minor status code.

### 4. IANA Considerations

This document has no IANA considerations currently. If and when a relevant IANA registry of GSS-API symbols and constants is created, then the GSS\_KRB5\_S\_KG\_INPUT\_TOO\_LONG minor status code should be added to such a registry.

## 5. Security Considerations

Kerberos V encryption types' PRF functions use a key derived from contexts' session keys and should preserve the forward security properties of the mechanisms' key exchanges.

Legacy Kerberos V encryption types may be weak, particularly the single-DES encryption types.

See also [RFC4401] for generic security considerations of GSS\_Pseudo\_random().

See also [RFC3961] for generic security considerations of the Kerberos V cryptographic framework.

Use of Ticket session keys, rather than sub-session keys, when initiators and acceptors fail to assert sub-session keys, is dangerous as ticket reuse can lead to key reuse; therefore, initiators should assert sub-session keys always, and acceptors should assert sub-session keys at least when initiators fail to do so.

The computational cost of computing this PRF+ may vary depending on the Kerberos V encryption types being used, but generally the computation of this PRF+ gets more expensive as the input and output octet string lengths grow (note that the use of a counter in the PRF+ construction allows for parallelization).

## 6. Normative References

- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", BCP 14, RFC 2119, DOI 10.17487/RFC2119, March 1997, <<http://www.rfc-editor.org/info/rfc2119>>.
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- [RFC4121] Zhu, L., Jaganathan, K., and S. Hartman, "The Kerberos Version 5 Generic Security Service Application Program Interface (GSS-API) Mechanism: Version 2", RFC 4121, DOI 10.17487/RFC4121, July 2005, <<http://www.rfc-editor.org/info/rfc4121>>.
- [RFC4401] Williams, N., "A Pseudo-Random Function (PRF) API Extension for the Generic Security Service Application Program Interface (GSS-API)", RFC 4401, DOI 10.17487/RFC4401, February 2006, <<http://www.rfc-editor.org/info/rfc4401>>.

## Appendix A. Test Vectors

Here are some test vectors from the MIT implementation provided by Greg Hudson. Test cases used include input string lengths of 0 and 61 bytes, and an output length of 44 bytes. 61 bytes of input is just enough to produce a partial second MD5 or SHA1 hash block with the four-byte counter prefix. 44 bytes of output requires two full and one partial RFC 3961 PRF output for all existing encyptes. All keys were randomly generated.

Encypte: des-cbc-crc  
Key: E607FE9DABB57AE0  
Input: (empty string)  
Output: 803C4121379FC4B87CE413B67707C4632EBED2C6D6B7  
2A55E878836E35E21600D915D590DED5B6D77BB30A1F

Encypte: des-cbc-crc  
Key: 54758316B6257A75  
Input: ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz123456789  
Output: 279E4105F7ADC9BD6EF28ABE31D89B442FE0058388BA  
33264ACB5729562DC637950F6BD144B654BE7700B2D6

Encypte: des3-cbc-sha1  
Key: 70378A19CD64134580C27C0115D6B34A1CF2FEECEF9886A2  
Input: (empty string)  
Output: 9F8D127C520BB826BFF3E0FE5EF352389C17E0C073D9  
AC4A333D644D21BA3EF24F4A886D143F85AC9F6377FB

Encypte: des3-cbc-sha1  
Key: 3452A167DF1094BA1089E0A20E9E51ABEF1525922558B69E  
Input: ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz123456789  
Output: 6BF24FABC858F8DD9752E4FCD331BB831F238B5BE190  
4EEA42E38F7A60C588F075C5C96A67E7F8B7BD0AECF4

Encypte: rc4-hmac  
Key: 3BB3AE288C12B3B9D06B208A4151B3B6  
Input: (empty string)  
Output: 9AEA11A3BCF3C53F1F91F5A0BA2132E2501ADF5F3C28  
3C8A983AB88757CE865A22132D6100EAD63E9E291AFA

Encypte: rc4-hmac  
Key: 6DB7B33A01BD2B72F7655CB7B3D5FA0B  
Input: ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz123456789  
Output: CDA9A544869FC84873B692663A82AFDA101C8611498B  
A46138B01E927C9B95EEC953B562807434037837DDDF

Entype: aes128-cts-hmac-sha1-96  
Key: 6C742096EB896230312B73972FA28B5D  
Input: (empty string)  
Output: 94208D982FC1BB7778128BDD77904420B45C9DA699F3  
117BCE66E39602128EF0296611A6D191A5828530F20F

Entype: aes128-cts-hmac-sha1-96  
Key: FA61138C109D834A477D24C7311BE6DA  
Input: ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz123456789  
Output: 0FAEDF0F842CC834FEE750487E1B622739286B975FE5  
B7F45AB053143C75CA0DF5D3D4BBB80F6A616C7C9027

Entype: aes256-cts-hmac-sha1-96  
Key: 08FCDAFD5832611B73BA7B497FEBFF8C954B4B58031CAD9B977C3B8C25192FD6  
Input: (empty string)  
Output: E627EFC14EF5B6D629F830C7109DEA0D3D7D36E8CD57  
A1F301C5452494A1928F05AFFBEE3360232209D3BE0D

Entype: aes256-cts-hmac-sha1-96  
Key: F5B68B7823D8944F33F41541B4E4D38C9B2934F8D16334A796645B066152B4BE  
Input: ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz123456789  
Output: 112F2B2D878590653CCC7DE278E9F0AA46FA5A380B62  
59F774CB7C134FCD37F61A50FD0D9F89BF8FE1A6B593

Entype: camellia128-cts-cmac  
Key: 866E0466A178279A32AC0BDA92B72AEB  
Input: (empty string)  
Output: 97FBB354BF341C3A160DCC86A7A910FDA824601DF677  
68797BACEEBF5D250AE929DEC9760772084267F50A54

Entype: camellia128-cts-cmac  
Key: D4893FD37DA1A211E12DD1E03E0F03B7  
Input: ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz123456789  
Output: 1DEE2FF126CA563A2A2326B9DD3F0095013257414C83  
FAD4398901013D55F367C82681186B7B2FE62F746BA4

Entype: camellia256-cts-cmac  
Key: 203071B1AE77BD3D6FCE70174AF95C225B1CED46B35CF52B6479EFEB47E6B063  
Input: (empty string)  
Output: 9B30020634C10FDA28420CEE7B96B70A90A771CED43A  
D8346554163E5949CBAE2FB8EF36AFB6B32CE75116A0

Entype: camellia256-cts-cmac  
Key: A171AD582C1AFBBAD52ABD622EE6B6A14D19BF95C6914B2BA40FFD99A88EC660  
Input: ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz123456789  
Output: A47CBB6E104DCC77E4DB48A7A474B977F2FB6A7A1AB6  
52317D50508AE72B7BE2E4E4BA24164E029CBACF786B

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### Authors' Addresses

Shawn Emery  
Oracle Corporation  
500 Eldorado Blvd Bldg 1  
Broomfield, CO 78727  
United States  
  
EMail: shawn.emery@oracle.com

Nicolas Williams  
Cryptonector, LLC  
  
EMail: nico@cryptonector.com