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## GOST R 34.11-2012: Hash Function

### Abstract

This document is intended to be a source of information about the Russian Federal standard hash function (GOST R 34.11-2012), which is one of the Russian cryptographic standard algorithms (called GOST algorithms). This document updates [RFC 5831](#).

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## 1. Scope

The Russian Federal standard hash function (GOST R 34.11-2012) establishes the hash-function algorithm and the hash-function calculation procedure for any sequence of binary symbols used in cryptographic methods of information processing and information security, including techniques for providing data integrity and authenticity and for digital signatures during information transfer, information processing, and information storage in computer-aided systems.

The hash function defined in the standard provides for the operation of digital signature systems using the asymmetric cryptographic algorithm in compliance with GOST R 34.10-2012 [[GOST3410-2012](#)].

GOST R 34.11-2012 applies to the creation, operation, and modernization of information systems of different purpose.

GOST R 34.11-94 is superseded by GOST R 34.11-2012 from 1st January 2013. That means that all new systems that are presented for certification MUST use GOST R 34.11-2012 and MAY use GOST R 34.11-94 also for maintaining compatibility with existing systems. Usage of GOST R 34.11-94 in current systems is allowed at least for a 5-year period.

This document updates [RFC 5831](#) [[RFC5831](#)].

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC 2119](#) [[RFC2119](#)].

## 2. General Information

1. GOST R 34.11-2012 [[GOST3411-2012](#)] was developed by the Center for Information Protection and Special Communications of the Federal Security Service of the Russian Federation with participation of the open joint-stock company Information Technologies and Communication Systems (InfoTeCS JSC).
2. GOST R 34.11-2012 was approved and introduced by Decree #216 of the Federal Agency on Technical Regulating and Metrology on 07.08.2012.
3. GOST R 34.11-2012 is intended to replace GOST R 34.11-94 [[GOST3411-94](#)], a national standard of the Russian Federation.

Terms and concepts in the standard comply with the following international standards:

- o ISO 2382-2 [[ISO2382-2](#)],
- o ISO/IEC 9796 [[ISO/IEC9796-2](#)][[ISO/IEC9796-3](#)],
- o series of standards ISO/IEC 14888 [[ISO/IEC14888-1](#)][[ISO/IEC14888-2](#)][[ISO/IEC14888-3](#)][[ISO/IEC14888-3Amd](#)], and
- o series of standards ISO/IEC 10118 [[ISO/IEC10118-1](#)][[ISO/IEC10118-2](#)][[ISO/IEC10118-3](#)][[ISO/IEC10118-4](#)].

## 3. Standard References

The following standards are referred to in GOST R 34.11-2012:

1. GOST 28147-89 [[GOST28147-89](#)], "Systems of information processing. Cryptographic data security. Algorithms of cryptographic transformation."

2. GOST R 34.10-2012 [[GOST3410-2012](#)], "Information technology. Cryptographic data security. Formation and verification processes of [electronic] digital signature."

Note: Users of the standard may check the validity of the referenced standards on the official Internet site of the Federal Agency on Technical Regulating and Metrology, in the annual reference book "National Standards" published on January 1 of the current year, and in corresponding monthly indices published during the current year. If the referenced standard is replaced (amended), then the replaced (amended) standard shall be used. If the referenced standard is canceled without replacement, then only the parts of this document not containing the specified reference may be used.

#### 4. Definitions and Notations

The following terms and their corresponding definitions are used in the standard.

##### 4.1. Definitions

padding: appending extra bits to a data string (Clause 3.9 of [ISO/IEC10118-1]).

initializing value: a value used in defining the starting point of a hash function (Clause 3.7 of [ISO/IEC10118-1]).

message: string of bits of any length (Clause 3.10 of [ISO/IEC14888-1]).

round function: a function that transforms two binary strings of lengths L1 and L2 to a binary string of length L2. It is used iteratively as part of a hash function, where it combines a data string of length L1 with the previous output of length L2 (Clause 3.10 of [ISO/IEC10118-1]).

Note: In GOST R 34.11-2012, the concepts "string of bits of length L" and "binary row vector of length L" are identical.

hash code: string of bits that is the output of a hash function (Clause 3.6 of [ISO/IEC14888-1]).

collision-resistant hash function: function that maps strings of bits to fixed-length strings of bits, satisfying the following properties:

1. for a given output, it is computationally infeasible to find an input that maps to this output;

2. for a given input, it is computationally infeasible to find a second input that maps to the same output; and
3. it is computationally infeasible to find any two distinct inputs that map to the same output (Clauses 3.2 and 3.7 of [ISO/IEC14888-1]).

Note: In the standard (to provide terminological compatibility with the current native standard documentation and with the published scientific and technical works), the terms "hash function" and "cryptographic hash function" are synonyms.

signature: one or more data elements resulting from the signature process (Clause 3.12 of [ISO/IEC 14888-1]).

Note: In the standard (to provide terminological compatibility with the current native standard documentation and with the published scientific and technical works), the terms "digital signature", "electronic signature", and "electronic digital signature" are synonyms.

#### 4.2. Notations

The following notations are used in the standard:

$V^*$	the set of all binary row vectors of finite length (hereinafter referred to as vectors) including empty string
$ A $	the length (number of components) of the vector $A$ belonging to $V^*$ (if $A$ is an empty string, then $ A  = 0$ )
$V_n$	the set of all binary vectors of length $n$ , where $n$ is a non-negative integer; subvectors and vector components are enumerated from right to left starting from zero
$(xor)$	exclusive-or of the two binary vectors of the same length
$A  B$	concatenation of vectors $A$ , $B$ (both belong to $V^*$ ), i.e., a vector from $V_{( A + B )}$ , where the left subvector from $V_{( A )}$ is equal to the vector $A$ and the right subvector from $V_{( B )}$ is equal to the vector $B$
$A^n$	concatenation of $n$ instances of the vector $A$
$Z_{(2^n)}$	ring of residues modulo $2^n$
$[+]$	addition operation in the ring $Z_{(2^n)}$

$\text{Vec}_n: \mathbb{Z}_2(2^n) \rightarrow V_n$   
 bijective-mapping operation associating an element from  $\mathbb{Z}_2(2^n)$  with its binary representation, i.e., for an element  $z$  of the ring  $\mathbb{Z}_2(2^n)$ , represented by the residue  $z_0 + (2 \cdot z_1) + \dots + (2^{n-1} \cdot z_{n-1})$ , where  $z_i \in \{0, 1\}$ ,  $i = 0, \dots, n-1$ , the equality  $\text{Vec}_n(z) = z_{n-1} || \dots || z_1 || z_0$  holds

$\text{Int}_n: V_n \rightarrow \mathbb{Z}_2(2^n)$   
 the mapping inverse to the mapping  $\text{Vec}_n$ , i.e.,  
 $\text{Int}_n = \text{Vec}_n^{-1}$

$\text{MSB}_n: V^* \rightarrow V_n$   
 the mapping associating the vector  $z_{k-1} || \dots || z_1 || z_0$ ,  $k \geq n$ , with the vector  $z_{k-1} || \dots || z_{k-n+1} || z_{k-n}$

$a := b$  operation of assigning the value  $b$  to the variable  $a$

$PS$  product of mappings, where the mapping  $S$  applies first

$M$  binary vector subject to hashing procedure,  $M$  belongs to  $V^*$ ,  
 $|M| < 2^{512}$

$H: V^* \rightarrow V_n$   
 hash function mapping the vector (message)  $M$  into the vector (hash code)  $H(M)$

$IV$  hash-function initializing value,  $IV \in V_{512}$

## 5. General Provisions

GOST R 34.11-2012 defines two hash functions  $H: V^* \rightarrow V_n$  with the hash-code lengths  $n = 512$  bits and  $n = 256$  bits.

## 6. Parameter Values

### 6.1. Initializing Values

The initializing value  $IV$  for a hash function with a hash-code length of 512 bits is  $0^{512}$ . The initializing value  $IV$  for a hash function with a hash-code length of 256 bits is  $(00000001)^{64}$ .

## 6.2. Nonlinear Bijections of Binary Vector Sets

Nonlinear bijection of the binary vector set  $V_8$  is presented by the following substitution:

$Pi = (Vec_8)Pi'(Int_8): V_8 \rightarrow V_8$

where  $Pi': Z_{(2^8)} \rightarrow Z_{(2^8)}$ .

The values of the substitution  $Pi'$  are presented in the array form  $Pi' = (Pi'(0), Pi'(1), \dots, Pi'(255))$ :

```
Pi' = (252, 238, 221, 17, 207, 110, 49, 22, 251, 196, 250,
      218, 35, 197, 4, 77, 233, 119, 240, 219, 147, 46,
      153, 186, 23, 54, 241, 187, 20, 205, 95, 193, 249,
      24, 101, 90, 226, 92, 239, 33, 129, 28, 60, 66,
      139, 1, 142, 79, 5, 132, 2, 174, 227, 106, 143,
      160, 6, 11, 237, 152, 127, 212, 211, 31, 235, 52,
      44, 81, 234, 200, 72, 171, 242, 42, 104, 162, 253,
      58, 206, 204, 181, 112, 14, 86, 8, 12, 118, 18,
      191, 114, 19, 71, 156, 183, 93, 135, 21, 161, 150,
      41, 16, 123, 154, 199, 243, 145, 120, 111, 157, 158,
      178, 177, 50, 117, 25, 61, 255, 53, 138, 126, 109,
      84, 198, 128, 195, 189, 13, 87, 223, 245, 36, 169,
      62, 168, 67, 201, 215, 121, 214, 246, 124, 34, 185,
      3, 224, 15, 236, 222, 122, 148, 176, 188, 220, 232,
      40, 80, 78, 51, 10, 74, 167, 151, 96, 115, 30,
      0, 98, 68, 26, 184, 56, 130, 100, 159, 38, 65,
      173, 69, 70, 146, 39, 94, 85, 47, 140, 163, 165,
      125, 105, 213, 149, 59, 7, 88, 179, 64, 134, 172,
      29, 247, 48, 55, 107, 228, 136, 217, 231, 137, 225,
      27, 131, 73, 76, 63, 248, 254, 141, 83, 170, 144,
      202, 216, 133, 97, 32, 113, 103, 164, 45, 43, 9,
      91, 203, 155, 37, 208, 190, 229, 108, 82, 89, 166,
      116, 210, 230, 244, 180, 192, 209, 102, 175, 194, 57,
      75, 99, 182)
```

### 6.3. Byte Permutation

The values of the permutation  $\tau$  belonging to  $S_{64}$  are presented in the array form  $\tau = (\tau(0), \tau(1), \dots, \tau(63))$ :

```

 $\tau = (0, 8, 16, 24, 32, 40, 48, 56,$ 
 $1, 9, 17, 25, 33, 41, 49, 57,$ 
 $2, 10, 18, 26, 34, 42, 50, 58,$ 
 $3, 11, 19, 27, 35, 43, 51, 59,$ 
 $4, 12, 20, 28, 36, 44, 52, 60,$ 
 $5, 13, 21, 29, 37, 45, 53, 61,$ 
 $6, 14, 22, 30, 38, 46, 54, 62,$ 
 $7, 15, 23, 31, 39, 47, 55, 63)$ 

```

### 6.4. Linear Transformations of Binary Vector Sets

Linear transformation  $l$  of the binary vector set  $V_{64}$  is specified by the right multiplication with the matrix  $A$  over the field  $GF(2)$ . The matrix rows are specified sequentially in a hexadecimal form. The row with number  $j$ ,  $j = 0, \dots, 63$  (specified in the form  $a_{\cdot}(j, 15) \dots a_{\cdot}(j, 0)$ , where  $a_{\cdot}(j, i)$  belongs to  $Z_{16}$ ,  $i = 0, \dots, 15$ ), is  $Vec_4(a_{\cdot}(j, 15)) || \dots || Vec_4(a_{\cdot}(j, 0))$ .

```

8e20faa72ba0b470 47107ddd9b505a38 ad08b0e0c3282d1c d8045870ef14980e
6c022c38f90a4c07 3601161cf205268d 1b8e0b0e798c13c8 83478b07b2468764
a011d380818e8f40 5086e740ce47c920 2843fd2067adea10 14aff010bdd87508
0ad97808d06cb404 05e23c0468365a02 8c711e02341b2d01 46b60f011a83988e
90dab52a387ae76f 486dd4151c3dfdb9 24b86a840e90f0d2 125c354207487869
092e94218d243cba 8a174a9ec8121e5d 4585254f64090fa0 accc9ca9328a8950
9d4df05d5f661451 c0a878a0a1330aa6 60543c50de970553 302a1e286fc58ca7
18150f14b9ec46dd 0c84890ad27623e0 0642ca05693b9f70 0321658cba93c138
86275df09ce8aaa8 439da0784e745554 afc0503c273aa42a d960281e9d1d5215
e230140fc0802984 71180a8960409a42 b60c05ca30204d21 5b068c651810a89e
456c34887a3805b9 ac361a443d1c8cd2 561b0d22900e4669 2b838811480723ba
9bcf4486248d9f5d c3e9224312c8c1a0 effa11af0964ee50 f97d86d98a327728
e4fa2054a80b329c 727d102a548b194e 39b008152acb8227 9258048415eb419d
492c024284fbaec0 aa16012142f35760 550b8e9e21f7a530 a48b474f9ef5dc18
70a6a56e2440598e 3853dc371220a247 1ca76e95091051ad 0edd37c48a08a6d8
07e095624504536c 8d70c431ac02a736 c83862965601dd1b 641c314b2b8ee083

```

Here one string contains 4 rows of the matrix  $A$ . So, the string with number  $i$ ,  $i = 0, \dots, 15$ , specifies 4 rows of the matrix  $A$  (with the numbers  $4i + j$ ,  $j = 0, \dots, 3$ ) in the following left-to-right order:  $4i + 0, 4i + 1, 4i + 2, 4i + 3$ .



The product of the vector  $b = b_{63} \dots b_0$  belonging to  $V_{64}$  and the matrix  $A$  is the vector  $c$  belonging to  $V_{64}$ :

$$\begin{aligned} c &= b_{63}(\text{Vec}_4(a_{(0, 15)})) \parallel \dots \parallel \text{Vec}_4(a_{(0, 0)}) && (\text{xor}) \\ &\dots && (\text{xor}) \\ &b_0(\text{Vec}_4(a_{(63, 15)})) \parallel \dots \parallel \text{Vec}_4(a_{(63, 0)}) \end{aligned}$$

where

$$\begin{aligned} b_i(\text{Vec}_4(a_{(63-i, 15)})) \parallel \dots \parallel \text{Vec}_4(a_{(63-i, 0)}) &= \\ &= 0^{64}, \text{ if } b_i = 0 \\ &= (\text{Vec}_4(a_{(63-i, 15)})) \parallel \dots \parallel \text{Vec}_4(a_{(63-i, 0)}), \text{ if } b_i = 1 \end{aligned}$$

for all  $i = 0, \dots, 63$ .

### 6.5. Iteration Constants

Iteration constants are specified in a hexadecimal form. The constant value specified in the form  $a_{127} \dots a_0$  (where  $a_i$  belongs to  $Z_{16}$ ,  $i = 0, \dots, 127$ ) is  $\text{Vec}_4(a_{127}) \parallel \dots \parallel \text{Vec}_4(a_0)$ :

$C[1] =$  b1085bdalecadae9ebcb2f81c0657c1f  
2f6a76432e45d016714eb88d7585c4fc  
4b7ce09192676901a2422a08a460d315  
05767436cc744d23dd806559f2a64507

$C[2] =$  6fa3b58aa99d2f1a4fe39d460f70b5d7  
f3feea720a232b9861d55e0f16b50131  
9ab5176b12d699585cb561c2db0aa7ca  
55dda21bd7cbcd56e679047021b19bb7

$C[3] =$  f574dcac2bce2fc70a39fc286a3d8435  
06f15e5f529c1f8bf2ea7514b1297b7b  
d3e20fe490359eb1c1c93a376062db09  
c2b6f443867adb31991e96f50aba0ab2

$C[4] =$  ef1fdfb3e81566d2f948e1a05d71e4dd  
488e857e335c3c7d9d721cad685e353f  
a9d72c82ed03d675d8b71333935203be  
3453eaa193e837f1220cbebc84e3d12e

$C[5] =$  4bea6bacad4747999a3f410c6ca92363  
7f151c1f1686104a359e35d7800fffb  
bfcd1747253af5a3dfff00b723271a16  
7a56a27ea9ea63f5601758fd7c6cfe57

```
C[6] = ae4faeae1d3ad3d96fa4c33b7a3039c0
      2d66c4f95142a46c187f9ab49af08ec6
      cffaa6b71c9ab7b40af21f66c2bec6b6
      bf71c57236904f35fa68407a46647d6e

C[7] = f4c70e16eeaac5ec51ac86febf240954
      399ec6c7e6bf87c9d3473e33197a93c9
      0992abc52d822c3706476983284a0504
      3517454ca23c4af38886564d3a14d493

C[8] = 9b1f5b424d93c9a703e7aa020c6e4141
      4eb7f8719c36de1e89b4443b4ddbc49a
      f4892bcb929b069069d18d2bd1a5c42f
      36acc2355951a8d9a47f0dd4bf02e71e

C[9] = 378f5a541631229b944c9ad8ec165fde
      3a7d3a1b258942243cd955b7e00d0984
      800a440bdbb2ceb17b2b8a9aa6079c54
      0e38dc92cb1f2a607261445183235adb

C[10] = abbedea680056f52382ae548b2e4f3f3
      8941e71cff8a78db1fffel8a1b336103
      9fe76702af69334b7a1e6c303b7652f4
      3698fad1153bb6c374b4c7fb98459ced

C[11] = 7bcd9ed0efc889fb3002c6cd635afe94
      d8fa6bbbebab07612001802114846679
      8a1d71efea48b9caefbacd1d7d476e98
      dea2594ac06fd85d6bcaa4cd81f32d1b

C[12] = 378ee767f11631bad21380b00449b17a
      cda43c32bcdfl1d77f82012d430219f9b
      5d80ef9d1891cc86e71da4aa88e12852
      faf417d5d9b21b9948bc924af11bd720
```

## 7. Transformations

For calculating the hash code  $H(M)$  of the message  $M$  belonging to  $V^*$ , the following transformations are used:

```
X[k]: V_512 -> V_512,
X[k](a) = k (xor) a, k, a belongs to V_512;
```

```
S:V_512 -> V_512,
S(a) = S(a_63||...||a_0) = Pi(a_63)||...||Pi(a_0), where
a = a_63||...||a_0 belongs to V_512, a_i belongs to V_8,
i = 0, ..., 63;
```

$P:V_{512} \rightarrow V_{512}$ ,  
 $P(a) = P(a_{63} || \dots || a_0) = a_{\text{Tau}(63)} || \dots || a_{\text{Tau}(0)}$ , where  
 $a = a_{63} || \dots || a_0$  belongs to  $V_{512}$ ,  $a_i$  belongs to  $V_8$ ,  
 $i = 0, \dots, 63$ ;  
  
 $L:V_{512} \rightarrow V_{512}$ ,  
 $L(a) = L(a_7 || \dots || a_0) = l(a_7) || \dots || l(a_0)$ , where  
 $a = a_7 || \dots || a_0$  belongs to  $V_{512}$ ,  $a_i$  belongs to  $V_{64}$ ,  
 $i = 0, \dots, 7$ .

## 8. Round Functions

The hash-code value of the message  $M$  belonging to  $V^*$  is calculated using the iterative procedure. Each iteration is provided using the round function:

$g_N:V_{512} \times V_{512} \rightarrow V_{512}$ , where  $N$  belongs to  $V_{512}$

calculated as

$g_N(h, m) = E(LPS(h \text{ (xor) } N), m) \text{ (xor) } h \text{ (xor) } m$

where

$E(K, m) = X[K[13]]LPSX[K[12]] \dots LPSX[K[2]]LPSX[K[1]](m)$

Values  $K[i]$  belonging to  $V_{512}$ ,  $i = 1, \dots, 13$ , are calculated as follows:

$K[1] = K$

$K[i] = LPS(K[i-1] \text{ (xor) } C[i-1]), i = 2, \dots, 13$

## 9. Hash-Function Calculation Procedure

Initial data for the procedure of calculating the hash code  $H(M)$  are a message  $M$  belonging to  $V^*$  (subject to hashing) and initializing value  $IV$  belonging to  $V_{512}$ . The algorithm for calculating the function  $H$  consists of the following steps.

Step 1. Assign initial values to the following variables:

1.1.  $h := IV$

1.2.  $N := 0^{512}$  belonging to  $V_{512}$

1.3.  $\text{EPSILON} := 0^{512}$  belonging to  $V_{512}$

1.4. Go to Step 2.

Step 2.

2.1. Check the condition  $|M| < 512$

If it is true, then go to Step 3.

Else, perform the following calculations:

2.2. Calculate the subvector  $m$  belonging to  $V_{512}$  of the message  $M$ :

$$M = M' || m$$

Then perform the following calculations:

2.3.  $h := g_N(h, m)$

2.4.  $N := \text{Vec}_{512}(\text{Int}_{512}(N) [+] 512)$

2.5.  $\text{EPSILON} := \text{Vec}_{512}(\text{Int}_{512}(\text{EPSILON}) [+] \text{Int}_{512}(m))$

2.6.  $M := M'$

2.7. Go to Step 2.1.

Step 3.

3.1.  $m := 0^{511-|M|} || 1 || M$

3.2.  $h := g_N(h, m)$

3.3.  $N := \text{Vec}_{512}(\text{Int}_{512}(N) [+] |M|)$

3.4.  $\text{EPSILON} := \text{Vec}_{512}(\text{Int}_{512}(\text{EPSILON}) [+] \text{Int}_{512}(m))$

3.5.  $h := g_0(h, N)$

3.6.  $h := g_0(h, \text{EPSILON})$ , for function with 512-bit hash code

$h := \text{MSB}_{256}(g_0(h, \text{EPSILON}))$ , for function with 256-bit hash code

3.7. End of the algorithm

The value of the variable  $h$  (obtained in Step 3.6) is the value of hash function  $H(M)$ .

## 10. Examples (Informative)

This section is for information only and is not a normative part of the standard.

The vectors from  $V^*$  are specified in a hexadecimal form. The vector  $A$  belonging to  $V_{(4n)}$  (specified in the form  $a_{(n-1)} \dots a_0$ , where  $a_i$  belongs to  $Z_{16}$ ,  $i = 0, \dots, n-1$ ) is  $\text{Vec}_4(a_{(n-1)}) || \dots || \text{Vec}_4(a_0)$ .

### 10.1. Example 1

Let's calculate the hash code of the following message (represented as a hexadecimal string):

```
M1 = 32313039383736353433323130393837
      36353433323130393837363534333231
      30393837363534333231303938373635
      343332313039383736353433323130
```

#### 10.1.1.1. For Hash Function with 512-Bit Hash Code

Assign the following values to the variables:

$$h := IV = 0^{512}$$
$$N := 0^{512}$$
$$\text{EPSILON} := 0^{512}$$

The length of the message is  $|M_1| = 504 < 512$ , so the incomplete block is padded:

```
m := 01323130393837363534333231303938
37363534333231303938373635343332
31303938373635343332313039383736
35343332313039383736353433323130
```

Calculate

$$K := \text{LPS}(h \text{ (xor) } N) = \text{LPS}(0^{512}).$$

After the transformation S:

[illegible]



$S(K[1] \text{ (xor) } C[1]) =$  ddf644e6e15f5733bff249410445536f  
 4e9bd69e200f3596b3d9ea737d70ald7  
 d1b6143b9c9288357758f8ef78278aa1  
 55f4d717dda7cb12b211e87e7f19203d

$PS(K[1] \text{ (xor) } C[1]) =$  ddbf4eb3d17755b2f6f29bd9b658f411  
 4449d6ea14f8d7e8e6419e733bef177e  
 e104207d9c78dd7f5f450f709227a719  
 575335a1888acb20336f96d735a1123d

$LPS(K[1] \text{ (xor) } C[1]) =$  d0b00807642fd78f13f2c3ebc774e80d  
 e0e902d23aef2ee9a73d010807dae9c1  
 88be14f0b2da27973569cd2ba0513010  
 36f728bd1d7eec33f4d18af70c46cfl e

Iteration 2

$K[2] =$  d0b00807642fd78f13f2c3ebc774e80d  
 e0e902d23aef2ee9a73d010807dae9c1  
 88be14f0b2da27973569cd2ba0513010  
 36f728bd1d7eec33f4d18af70c46cfl e

$LPSX[K[2]]LPSX[K[1]](m) =$  18e77571e703d19548075c574ce5e50e  
 0480c9c5b9f21d45611ab86cf32e352a  
 d91854ea7df8f863d46333673f62ff2d  
 3efaelcd966f8e2a74ce49902799aad4

Iteration 3

$K[3] =$  9d4475c7899f2d0bb0e8b7dac6ef6e6b  
 44ecf66716d3a0f16681105e2d13712a  
 1a9387ecc257930e2d61014a1b5c9fc9  
 e24e7d636eb1607e816dbaf927b8fca9

$LPSX[K[3]] \dots LPSX[K[1]](m) =$  03dc0a9c64d42543ccdb62960d58c17e  
 0b5b805d08a07406ece679d5f82b70fe  
 a22a7ea56e21814619e8749b30821457  
 5489d4d465539852cd4b0cd3829bef39

Iteration 4

$K[4] =$  5c283daba5ec1f233b8c833c48e1c670  
 dae2e40cc4c3219c73e58856bd96a72f  
 df9f8055ffe3c004c8cde3b8bf78f95f  
 3370d0a3d6194ac5782487defd83ca0f

LPSX[K[4]]...LPSX[K[1]](m) = dbee312ea7301b0d6d13e43855e85db8  
1608c780c43675bc93cfd82c1b4933b3  
898a35b13e1878abe119e4dfffb9de488  
9738ca74d064cd9eb732078c1fb25e04

Iteration 5

K[5] = 109f33262731f9bd569cbc9317baa551  
d4d2964fa18d42c41fab4e37225292ec  
2fd97d7493784779046388469ae195c4  
36fa7cba93f8239ceb5ffc818826470c

LPSX[K[5]]...LPSX[K[1]](m) = 7fb3f15718d90e889f9fb7c38f527bec  
861c298afb9186934a93c9d96ade20df  
109379bb9c1a1ffd0ad81fce7b45ccd5  
4501e7d127e32874b5d7927b032de7a1

Iteration 6

K[6] = b32c9b02667911cf8f8a0877be9a1707  
57e25026ccf41e67c6b5da70b1b87474  
3e1135cfbefe244237555c676c153d99  
459bc382573aee2d85d30d99f286c5e7

LPSX[K[6]]...LPSX[K[1]](m) = 95efa4e104f235824bae5030fe2d0f17  
0a38de3c9b8fc6d8fala9adc2945c413  
389a121501fa71a65067916b0c06f6b8  
7ce18de1a2a98e0a64670985f47d73f1

Iteration 7

K[7] = 8a13c1b195fd0886ac49989e7d84b08b  
c7b00e4f3f62765ece6050fcbabdc234  
6c8207594714e8e9c9c7aad694edc922  
d6b01e17285eb7e61502e634559e32f1

LPSX[K[7]]...LPSX[K[1]](m) = 7ea4385f7e5e40103bfb25c67e404c75  
24eec43e33b1d06557469c6049854304  
32b43d941b77ffd476103338e9bd5145  
d9c1e18b1f262b58a81dcefff6fc6535

Iteration 8

K[8] = 52cec3b11448bb8617d0ddfb9c926f2e8  
8730cb9179d6decea5acbfd323ec376  
4c47f7a9e13bb1db56c342034773023d  
617ff01cc546728e71dff8de5d128cac



LPSX[K[8]]...LPSX[K[1]](m) = b2426da0e58d5cfe898c36e797993f90  
2531579d8ecc59f8dd8a60802241a456  
1f290cf992eb398894424bf681636968  
c167e870967b1dd9047293331956daba

Iteration 9

K[9] = f38c5b7947e7736d502007a05ea64a4e  
b9c243cb82154aa138b963bbb7f28e74  
d4d710445389671291d70103f48fd4d4  
c01fc415e3fb7dc61c6088afalale735

LPSX[K[9]]...LPSX[K[1]](m) = 5e0c9978670b25912dd1ede5bdd1cf18  
ed094d14c6d973b731d50570d0a9bca2  
15415a15031fd20ddefb5bc61b96671d  
6902f49df4d2fd346ceebda9431cb075

Iteration 10

K[10] = 0740b3faa03ed39b257dd6e3db7c1bf5  
6b6e18e40cdaabd30617cecbaddd618e  
a5e61bb4654599581dd30c24clab877a  
d0687948286cfefaa7eef99f6068b315

LPSX[K[10]]...LPSX[K[1]](m) = c1ddd840fe491393a5d460440e03bf45  
1794e792c0c629e49ab0c1001782dd37  
691cb6896f3e00b87f71d37a584c35b9  
cd8789fad55a46887e5b60e124b51a61

Iteration 11

K[11] = 185811cf3c2633aec8cfdcfcae9dbb293  
47011bf92b95910a3ad71e5fca678e45  
e374f088f2e5c29496e9695ce8957837  
107bb3aa56441af11a82164893313116

LPSX[K[11]]...LPSX[K[1]](m) = 3f75beaf2911c35d575088e30542b689  
c85b6b1607f8b800405941f5ab704284  
7b9b08b58b4fbdd6154ed7b366fd3ee7  
78ce647726ddb3c7d48c8ce8866a8435

Iteration 12

K[12] = 9d46bf66234a7ed06c3b2120d2a3f15e  
0fedd87189b75b3cd2f206906b5ee00d  
c9aleab800fb8cc5760b251f4db5cdef  
427052fa345613fd076451901279ee4c

```
LPSX[K[12]]...LPSX[K[1]](m) = f35b0d889eadfcff73b6b17f33413a97
                                417d96f0c4cc9d30cda8ebb7dcd5d1b0
                                61e620bac75b367370605f474ddc0060
                                03bec4c4d7ce59a73fbe6766934c55a2
```

Iteration 13

```
K[13] = 0f79104026b900d8d768b6e223484c97
        61e3c585b3a405a6d2d8565ada926c3f
        7782ef127cd6b98290bf612558b4b60a
        a3cbc28fd94f95460d76b621cb45be70
```

```
X[K[13]]...LPSX[K[1]](m) = fc221dc8b814fc27a4de079d10097600
                            209e5375776898961f70bded0647bd8f
                            1664cfa8bb8d8ff1e0df3e621568b66a
                            a075064b0e81cce132c8d1475809ebd2
```

The result of the transformation  $g_N(h, m)$  is

```
h = fd102cf8812ccb1191ea34af21394f38
    17a86641445aa9a626488adb33738ebd
    2754f6908cbbbac5d3ed0f522c50815c
    954135793fblf5d905fee4736b3bdae2
```

Variables  $N$  and  $\text{EPSILON}$  change their values to

```
N = 00000000000000000000000000000000
    00000000000000000000000000000000
    00000000000000000000000000000000
    000000000000000000000000000000001f8
```

```
EPSILON = 01323130393837363534333231303938
           37363534333231303938373635343332
           31303938373635343332313039383736
           35343332313039383736353433323130
```

The result of the transformation  $g_0(h, N)$  is

```
h = 5c881fd924695cf196c2e4fec20d14b6
    42026f2a0b1716ebaabb7067d4d59752
    3d2db69d6d3794622147a14f19a66e7f
    9037e1d662d34501a8901a5de7771d7c
```

The result of the transformation  $g_0(h, \text{EPSILON})$  is

```
h = 486f64c1917879417fef082b3381a4e2
    11c324f074654c38823a7b76f830ad00
    fa1fbae42b1285c0352f227524bc9ab1
    6254288dd6863dccd5b9f54alad0541b
```

The hash code of the message  $M_1$  is the value

```
H(M1) = 486f64c1917879417fef082b3381a4e2
        11c324f074654c38823a7b76f830ad00
        fa1fbae42b1285c0352f227524bc9ab1
        6254288dd6863dccd5b9f54alad0541b
```

#### 10.1.2. For Hash Function with 256-Bit Hash Code

Assign the following values to the variables:

```
h := IV = (00000001)^64
```

```
N := 0^512
```

```
EPSILON := 0^512
```

The length of the message is  $|M_1| = 504 < 512$ , so the incomplete block is padded:

```
m := 01323130393837363534333231303938
      37363534333231303938373635343332
      31303938373635343332313039383736
      35343332313039383736353433323130
```

Calculate

```
K := LPS(h (xor) N) = LPS((00000001)^64)
```

After the transformation  $S$ :

```
S(h (xor) N) = eeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeee
                eeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeee
                eeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeee
                eeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeee
```

after the transformation P:

[illegible]

after the transformation L:

```
K := LPS(h (xor) N) = 23c5ee40b07b5f1523c5ee40b07b5f15
                        23c5ee40b07b5f1523c5ee40b07b5f15
                        23c5ee40b07b5f1523c5ee40b07b5f15
                        23c5ee40b07b5f1523c5ee40b07b5f15
```

Then the transformation  $E(K, m)$  is performed:

Iteration 1

$$K[1] = \begin{matrix} 23c5ee40b07b5f1523c5ee40b07b5f15 \\ 23c5ee40b07b5f1523c5ee40b07b5f15 \\ 23c5ee40b07b5f1523c5ee40b07b5f15 \\ 23c5ee40b07b5f1523c5ee40b07b5f15 \end{matrix}$$

```
X[K[1]](m) = 22f7df708943682316f1dd72814b662d
               14f3db7483496e251afdd976854f6c27
               12f5d778874d6a2110f7df7089436823
               16f1dd72814b662d14f3db7483496e25
```

```
SX[K[1]](m) = 65c061327951f35a99a6d819f5a29a01
                93d290ffa92ab25cf14b538aa8cc9d21
                f0f4fe6dc93a7818e9c061327951f35a
                99a6d819f5a29a0193d290ffa92ab25c
```

```
PSX[K[1]](m) = 659993f1f0e99993c0a6d24bf4c0a6d2
61d89053fe61d8903219ff8a6d3219ff
79f5a9a8c979f5a951a22acc3a51a22a
f39ab29d78f39ab25a015c21185a015c
```

```
LPSX[K[1]](m) = e549368917a0a2611d5e08c9c2fd5b3c
563f18c0f68c410d84ae9d5fbdfb9340
55650121b7aa6d7b3e7d09d46ac4358a
daa6ae44fa3b0402c4166d2c3eb2ef02
```

```
K[1] (xor) C[1] = 92cdb59aaeb185fcc80ec1c1701e230a
0caf98039e3e8f03528b56cdc5fe9be9
68b90ed1221c36148187c448141b8c00
26b39a767c0f1236fe458b1942dd1a12
```

$S(K[1] \text{ (xor) } C[1]) =$  ecd95e282645a83930045858325f5afa  
2341dc110ad303110ef676d9ac63509b  
f3a3041b65148f93f5c986f293bb7cfc  
ef92288ac34df08f63c8f6362cd8f1f0

$PS(K[1] \text{ (xor) } C[1]) =$  ec30230ef3f5ef63d90441f6a3c992c8  
5e58dc76048628f6285811d91bf28a36  
26320aac6593c32c455fd36314bb4dd8  
a85a03508f7cf0f139fa119b93fc8ff0

$LPS(K[1] \text{ (xor) } C[1]) =$  18ee8f3176b2ebea3bd6cb8233694cea  
349769df88be26bf451cfab6a904a549  
da22de93a66a66b19c7e6b5eea633511  
e611d68c8401bfcd0c7d0cc39d4a5eb9

Iteration 2

$K[2] =$  18ee8f3176b2ebea3bd6cb8233694cea  
349769df88be26bf451cfab6a904a549  
da22de93a66a66b19c7e6b5eea633511  
e611d68c8401bfcd0c7d0cc39d4a5eb9

$LPSX[K[2]]LPSX[K[1]](m) =$  c502dab7e79eb94013fcd1ba64def3b9  
16f18b63855d43d22b77fca1452f9866  
c2b45089c62e9d82edf1ef45230db9a2  
3c9e1c521113376628a5f6a5dbc041b2

Iteration 3

$K[3] =$  aaa4cf31a265959157aec8ce91e7fd46  
bf27dee21164c5e3940bba1a519e9d1f  
ce0913f1253e7757915000cd674be12c  
c7f68e73ba26fb00fd74af4101805f2d

$LPSX[K[3]] \dots LPSX[K[1]](m) =$  8e5a4fe41fc790af29944f027aa2f101  
05d65cf60a66e442832bb9ab5020dc54  
772e36b03d4b9aa471037212cde93375  
226552392ef4d83010a007e1117a07b5

Iteration 4

$K[4] =$  61fe0a65cc177af50235e2afadded326  
a5329a2236747bf8a54228aeca9c4585  
cd801ea9dd743a0d98d01ef0602b0e33  
2067fb5ddd6ac1568200311920839286

LPSX[K[4]]...LPSX[K[1]](m) = deeb40df69997afef726f03bdc13cb6  
ba9287698201296f2fd8284f06d33ea4  
a850a0ff48026dd47c1e88ec813ed2eb  
1186059d842d8d17f0bfa259e56655b1

Iteration 5

K[5] = 9983685f4fd3636f1fd5abb75fbf26a8  
e2934314aa2ecb3ee4693c86c06c7d4e  
169bd540af75e1610a546acd63d960ba  
d595394cc199bf6999a5d5309fe73d5a

LPSX[K[5]]...LPSX[K[1]](m) = 675ea894d326432e1af7b201bc369f8a  
b021f6fa58da09678ffc08ef30db43a3  
7f1f7347cb77da0f6ba30c85848896c3  
bac240ab14144283518b89a33d0caf07

Iteration 6

K[6] = f05772ae2ce7f025156c9a7fbcc6b8fd  
f1e735d613946e32922994e52820ffea  
62615d907eb0551ad170990a86602088  
af98c83c22cdb0e2be297c13c0f7a156

LPSX[K[6]]...LPSX[K[1]](m) = 1bc204bf9506ee9b86bbcf82d254a112  
aea6910b6db3805e399cb718d1b33199  
64459516967cee4e648e8cfbf81f56dc  
8da6811c469091be5123e6a1d5e28c73

Iteration 7

K[7] = 5ad144c362546e4e46b3e7688829fbb7  
7453e9c3211974330b2b8d0e6be2b5ac  
c89eb6b35167f159b7b005a43e5959a6  
51a9b18cfc8e4098fcf03d9b81cfbb8d

LPSX[K[7]]...LPSX[K[1]](m) = f30d791ed78bdee819022a3d78182242  
124efcdd54e203f23fb2dc7f94338ff9  
55a5afc15ffef03165263c4fdb36933a  
a982016471fbac9419f892551e9e568b

Iteration 8

K[8] = 6a6cec9a1ba20a8db64fa840b934352b  
518c638ed530122a83332fe0b8efdac9  
018287e5a9f509c78d6c746adcd5426f  
b0a0ad5790dfb73fc1f191a539016daa

LPSX[K[8]]...LPSX[K[1]](m) = 1fc20f1e91a1801a4293d3f3aa9e9156  
0fcc3810bb15f3ee9741c9b87452519f  
67cb9145519884a24de6db736a5cb143  
0da7458e5e51b80be5204ba5b2600177

Iteration 9

K[9] = 99217036737aa9b38a8d6643f705bd51  
f351531f948f0fc5e35fa35fee9dd8bd  
bb4c9d580a224e9cd82e0e2069fc49ed  
367d5f94374435382b8fb6a8f5dd0409

LPSX[K[9]]...LPSX[K[1]](m) = 1a52f09d1e81515a36171e0b1a2809c5  
0359bed90f2e78cbd89b7d4afa6d0466  
55c96bdae6ee97055cc7e857267c2ccf  
28c8f5dd95ed58a9a68c12663bb28967

Iteration 10

K[10] = 906763c0fc89fala69288d8ec9e9dda  
9a7630e8bfd6c3fed703c35d2e62aeaf  
f0b35d80a7317a7f76f83022f2526791  
ca8fdf678fcb337bd74fe5393ccb05d2

LPSX[K[10]]...LPSX[K[1]](m) = 764043744a0a93687e65aba8cfc25ec8  
714fb8e1bdc9ae2271e7205eaaa577c1  
b3b83e7325e50a19bd2d56b061b5de39  
235c9c9fd95e071a1a291a5f24e8c774

Iteration 11

K[11] = 88ce996c63618e6404a5c8e03ee43385  
4e2ae3eee68991bbbff3c29d38dad6e  
d6aldae9a6dc6ddf52ce34af272f96d3  
159c8c624c3fe6e13d695c0bfc89add5

LPSX[K[11]]...LPSX[K[1]](m) = 9b1ce8ff26b445cb288c0aeccf84658e  
ea91dbdf14828bf70110a5c9bd146cd9  
646350cff4e90e7b63c5cc325e9b4410  
81935f282d4648d9584f71860538f03b

Iteration 12

K[12] = 3e0a281ea9bd46063eec550100576f3a  
506aa168cf82915776b978fcca32f38  
b55f30c79982ca45628e8365d8798477  
e75a49c68199112ald7b5a0f7655f2db

```
LPSX[K[12]]...LPSX[K[1]](m) = 133aeecede251eb81914b8ba48dcbc0b
                                8a6fc63a292cc49043c3d3346b3f0829
                                a9cb71ecff25ed2a91bdcf8f649907c1
                                10cb76ff2e43100cdd4ba8a147a572f5
```

Iteration 13

```
K[13] = f0b273409eb31aebe432fbae18672122
        62c848422b6a92f93f6cbab54ed18b83
        14b21cffc51e3fa319ff433e76ef6adb
        0ef9f5e03c907falcf9eca06500bf03
```

```
X[K[13]]...LPSX[K[1]](m) = e3889d8e40960453fd26431450bb9d29
                             e8a78e78024656697caf698125ee83aa
                             bd796d133a3bd28988428cb112766d1a
                             1e32831f12d36fad21b2440122a5cdf6
```

The result of the transformation  $g_N(h, m)$  is

```
h = e3bbadbf78af3264c9137127608aa510
    de90ba4d3075665844965fb611dbb199
    8d48552a0c0ce6bcba71bc802a4f5b2d
    2a07b12c22e25794178570341096fdc7
```

The variables  $N$  and  $\text{EPSILON}$  change their values to

```
N = 00000000000000000000000000000000
    00000000000000000000000000000000
    00000000000000000000000000000000
    000000000000000000000000000000001f8
```

```
EPSILON = 01323130393837363534333231303938
           37363534333231303938373635343332
           31303938373635343332313039383736
           35343332313039383736353433323130
```

The result of the transformation  $g_0(h, N)$  is

```
h = 70f22bada4cfe18a6a56ec4b3f328cd4
    0db8e1bf8a9d5f711d5efab11191279d
    715aab7648d07eddbf87dc79c80516e6
    ffcbcf5678b0ac29ea00fa85c8173cc6
```



The result of the transformation  $g_0(h, \text{EPSILON})$  is

```
h = 00557be5e584fd52a449b16b0251d05d
    27f94ab76cbaa6da890b59d8ef1e159d
    2088e482e2acf564e0e9795a51e4dd26
    1f3f667985a2fcc40ac8631facal709a
```

The hash code of the message  $M_1$  is the value

```
H(M1) = 00557be5e584fd52a449b16b0251d05d
        27f94ab76cbaa6da890b59d8ef1e159d
```

## 10.2. Example 2

Let's calculate the hash code of the following message:

```
M2 = fbe2e5f0eee3c820fbaeafabef20fffb
      f0e1e0f0f520e0ed20e8ece0ebe5f0f2
      f120fff0eeec20f120faf2fee5e2202c
      e8f6f3ede220e8e6eeele8f0f2d1202c
      e8f0f2e5e220e5d1
```

### 10.2.1. For Hash Function with 512-Bit Hash Code

Assign the following values to the variables:

$h := IV = 0^{512}$

$N := 0^{512}$

$\text{EPSILON} := 0^{512}$

The length of the message is  $|M_2| = 576 > 512$ , so a part of this message is initially transformed:

```
m := fbaeafabef20fffbf0e1e0f0f520e0ed
      20e8ece0ebe5f0f2f120fff0eeec20f1
      20faf2fee5e2202ce8f6f3ede220e8e6
      eeele8f0f2d1202ce8f0f2e5e220e5d1
```

Calculate

$K := \text{LPS}(h \text{ (xor) } N) = \text{LPS}(0^{512})$



```
LPSX[K[1]](m) = 909aa733e1f52321a2fe35bfb8f67e92
                fbc70ef544709d5739d8faaca4acf126
                e83e273745c25b7b8f4a83a7436f6353
                753cbbbe492262cd3a868eace0104af1
```

```
K[1] (xor) C[1] = 028ba7f4d01e7f9d5848d3af0eb1d96b
                9ce98a6de0917562c2cd44a3bb516188
                f8ff1cbf5cb3cc7511c1d6266ab47661
                b6f5881802a0e8576e0399773c72e073
```

```
S(K[1] (xor) C[1]) = ddf644e6e15f5733bff249410445536f
                    4e9bd69e200f3596b3d9ea737d70a1d7
                    dl1b6143b9c9288357758f8ef78278aa1
                    55f4d717dda7cb12b211e87e7f19203d
```

```
PS(K[1] (xor) C[1]) = ddbf4eb3d17755b2f6f29bd9b658f411
                    4449d6ea14f8d7e8e6419e733bef177e
                    e104207d9c78dd7f5f450f709227a719
                    575335a1888acb20336f96d735a1123d
```

```
LPS(K[1] (xor) C[1]) = d0b00807642fd78f13f2c3ebc774e80d
                    e0e902d23aef2ee9a73d010807dae9c1
                    88be14f0b2da27973569cd2ba0513010
                    36f728bd1d7eec33f4d18af70c46cf1e
```

Iteration 2

```
K[2] = d0b00807642fd78f13f2c3ebc774e80d
        e0e902d23aef2ee9a73d010807dae9c1
        88be14f0b2da27973569cd2ba0513010
        36f728bd1d7eec33f4d18af70c46cf1e
```

```
LPSX[K[2]]LPSX[K[1]](m) = 301aadd761d13df0b473055b14a2f74a
                    45f408022aecadd4d5f19cab8228883a
                    021ac0b62600a495950c628354ffce11
                    61c68b7be7e0c58af090ce6b45e49f16
```

Iteration 3

```
K[3] = 9d4475c7899f2d0bb0e8b7dac6ef6e6b
        44ecf66716d3a0f16681105e2d13712a
        1a9387ecc257930e2d61014a1b5c9fc9
        e24e7d636eb1607e816dbaf927b8fca9
```

```
LPSX[K[3]]...LPSX[K[1]](m) = 9b83492b9860a93cbca1c0d8e0ce59db
                    04e10500a6ac85d4103304974e78d322
                    59ceff03fbb353147a9c948786582df7
                    8a34c9bde3f72b3ca41b9179c2cceef3
```

## Iteration 4

K[4] = 5c283daba5ec1f233b8c833c48e1c670  
dae2e40cc4c3219c73e58856bd96a72f  
df9f8055ffe3c004c8cde3b8bf78f95f  
3370d0a3d6194ac5782487defd83ca0f

LPSX[K[4]]...LPSX[K[1]](m) = e638e0a1677cdea107ec3402f70698a4  
038450dab44ac7a447e10155aa33ef1b  
daf8f49da7b66f3e05815045fbd39c99  
1cb0dc536e09505fd62d3c2cd00b0f57

## Iteration 5

K[5] = 109f33262731f9bd569cbc9317baa551  
d4d2964fa18d42c41fab4e37225292ec  
2fd97d7493784779046388469ae195c4  
36fa7cba93f8239ceb5ffc818826470c

LPSX[K[5]]...LPSX[K[1]](m) = 1c7c8e19b2bf443eb3adc0c787a52a17  
3821a97bc5a8efea58fb8b27861829f6  
dd5ff9c97865e08clac66f47392b578e  
21266e323a0aacedeec3ef0314f517c6

## Iteration 6

K[6] = b32c9b02667911cf8f8a0877be9a1707  
57e25026ccf41e67c6b5da70b1b87474  
3e1135cfbefe244237555c676c153d99  
459bc382573aee2d85d30d99f286c5e7

LPSX[K[6]]...LPSX[K[1]](m) = 48fecfc5b3eb77998fb39bfcccd128cd  
42fccb714221be1e675a1c6fdde7e311  
98b318622412af7e999a3eff45e6d616  
09a7f2ae5c2ff1ab7ff3b37be7011ba2

## Iteration 7

K[7] = 8a13c1b195fd0886ac49989e7d84b08b  
c7b00e4f3f62765ece6050fcbabdc234  
6c8207594714e8e9c9c7aad694edc922  
d6b01e17285eb7e61502e634559e32f1

LPSX[K[7]]...LPSX[K[1]](m) = a48f8d781c2c5be417ae644cc2e15a9f  
01fcead3232e5bd53f18a5ab875ccelb  
8ala400cf48521c7ce27fb1e94452fb5  
4de23118f53b364ee633170a62f5a8a9

## Iteration 8

K[8] = 52cec3b11448bb8617d0ddfb9c926f2e8  
8730cb9179d6decea5acbfffd323ec376  
4c47f7a9e13bb1db56c342034773023d  
617ff01cc546728e71dff8de5d128cac

LPSX[K[8]]...LPSX[K[1]](m) = e8a31b2e34bd2ae21b0ecf29cc4c37c7  
5c4d11d9b82852517515c23e81e906a4  
51b72779c3087141f1a15ab57f96d7da  
6c7ee38ed25befbdef631216356ff59c

## Iteration 9

K[9] = f38c5b7947e7736d502007a05ea64a4e  
b9c243cb82154aa138b963bbb7f28e74  
d4d710445389671291d70103f48fd4d4  
c01fc415e3fb7dc61c6088afalale735

LPSX[K[9]]...LPSX[K[1]](m) = 34392ed32ea3756e32979cb0a2247c39  
18e0b38d6455ca88183356bf8e5877e5  
5d542278a696523a8036af0f1c2902e9  
cbc585de803ee4d26649c9e1f00bda31

## Iteration 10

K[10] = 0740b3faa03ed39b257dd6e3db7c1bf5  
6b6e18e40cdaabd30617cecbaddd618e  
a5e61bb4654599581dd30c24c1ab877a  
d0687948286cfefaa7eef99f6068b315

LPSX[K[10]]...LPSX[K[1]](m) = 6a82436950177fea74cce6d507a5a64e  
54e8a3181458e3bdfbdbc6180c9787de  
7ccb676dd809e7cb1eb2c9ebd0165615  
70801a4e9ce17a438b85212f4409bb5e

## Iteration 11

K[11] = 185811cf3c2633aec8cfdcfcae9dbb293  
47011bf92b95910a3ad71e5fca678e45  
e374f088f2e5c29496e9695ce8957837  
107bb3aa56441af11a82164893313116

LPSX[K[11]]...LPSX[K[1]](m) = 7b97603135e2842189b0c9667596e96b  
d70472ccbc73ae89da7d1599c72860c2  
85f5771088f1fb0f943d949f22f1413c  
991eafb51ab8e5ad8644770037765aec

Iteration 12

K[12] = 9d46bf66234a7ed06c3b2120d2a3f15e  
0fedd87189b75b3cd2f206906b5ee00d  
c9a1eab800fb8cc5760b251f4db5cdef  
427052fa345613fd076451901279ee4c

LPSX[K[12]]...LPSX[K[1]](m) = 39ec8a88db635b46c4321adf41fd9527  
a39a67f6d7510db5044f05efaf721db5  
cf976a726ef33dc4dfcda94033e741a4  
63770861a5b25fefcb07281eed629c0e

Iteration 13

K[13] = 0f79104026b900d8d768b6e223484c97  
61e3c585b3a405a6d2d8565ada926c3f  
7782ef127cd6b98290bf612558b4b60a  
a3cbc28fd94f95460d76b621cb45be70

X[K[13]]...LPSX[K[1]](m) = 36959ac8fdda5b9e135aac3d62b5d9b0  
c279a27364f50813d69753b575e0718a  
b8158560122584464f72c8656b53f7ae  
c0bccae7cfdcaa9c6719e3f2627227e

The result of the transformation  $g_N(h, m)$  is

h = cd7f602312faa465e3bb4ccd9795395d  
e2914e938f10f8e127b7ac459b0c517b  
98ef779ef7c7a46aa7843b8889731f48  
2e5d221e8e2cea852e816cdac407c7af

The variables N and EPSILON change their values to

N = 00000000000000000000000000000000  
00000000000000000000000000000000  
00000000000000000000000000000000  
00000000000000000000000000000200

EPSILON = fbeafaebef20ffffbf0e1e0f0f520e0ed  
20e8ece0ebe5f0f2f120fff0eeec20f1  
20faf2fee5e2202ce8f6f3ede220e8e6  
eee1e8f0f2d1202ce8f0f2e5e220e5d1

The length of the rest of the message is less than 512, so the incomplete block is padded:

[illegible]

The result of the transformation  $g_N(h, m)$  is

```
h = c544ae6efdf14404f089c72d5faf8dc6  
aca1db5e28577fc07818095f1df70661  
e8b84d0706811cf92dffbf8f96e61493d  
c382795c6ed7a17b64685902cbdc878e
```

The variables N and EPSILON change their values to

```
N      = 00000000000000000000000000000000  
         00000000000000000000000000000000  
         00000000000000000000000000000000  
         0000000000000000000000000000240
```

```
EPSILON = fbeafaebef20fffbf0e1e0f0f520e0ed
20e8ece0ebe5f0f2f120fff0eeec20f1
20faf2fee5e2202ce8f6f3ede220e8e6
eee1e8f0f2d1202ee4d3d8d6d104adf1
```

The result of the transformation  $g_0(h, N)$  is

## h = 4deb6649ffa5caf4163d9d3f9967fbbd6eb3da68f916b6a09f41f2518b81292b703dc5d74e1ace5bcd3458af43bb456e837326088f2b5df14bf83997a0b1ad8d

The result of the transformation  $g_0(h, \text{EPSILON})$  is

## h = 28fbc9bada033b1460642bdcddb90c3fb3e56c497ccd0f62b8a2ad4935e85f037613966de4ee00531ae60f3b5a47f8dae06915d5f2f194996fcabf2622e6881e

The hash code of the message M2 is the value

H(M2) = 28fbc9bada033b1460642bdcddb90c3f  
b3e56c497ccd0f62b8a2ad4935e85f03  
7613966de4ee00531ae60f3b5a47f8da  
e06915d5f2f194996fcabf2622e6881e

#### 10.2.2. For Hash Function with 256-Bit Hash Code

Assign the following values to the variables:

```
h := IV = (00000001)^64
```

$$N := 0^{512}$$

```
EPSILON := 0^512
```

The length of the message is  $|M_2| = 576 > 512$ , so a part of this message is initially transformed:

```
m := fbeafaebef20ffffbf0e1e0f0f520e0ed
20e8ece0ebe5f0f2f120fff0eeec20f1
20faf2fee5e2202ce8f6f3ede220e8e6
eee1e8f0f2d1202ce8f0f2e5e220e5d1
```

Calculate:

```
K := LPS(h (xor) N) = LPS((00000001)^64)
```

After the transformation S:

```
S(h (xor) N) =
```

after the transformation P:

```
PS(h (xor) N) = 
XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX
XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX
XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX
XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX
```

after the transformation L:

```
K := LPS(h (xor) N) = 23c5ee40b07b5f1523c5ee40b07b5f15
                        23c5ee40b07b5f1523c5ee40b07b5f15
                        23c5ee40b07b5f1523c5ee40b07b5f15
                        23c5ee40b07b5f1523c5ee40b07b5f15
```



Then the transformation  $E(K, m)$  is performed:

Iteration 1

$K[1]$  = 23c5ee40b07b5f1523c5ee40b07b5f15  
 23c5ee40b07b5f1523c5ee40b07b5f15  
 23c5ee40b07b5f1523c5ee40b07b5f15  
 23c5ee40b07b5f1523c5ee40b07b5f15

$X[K[1]](m)$  = d82f14ab5f5ba0eed3240eb0455bbff8  
 032d02a05b9eafe7d2e511b05e977fe4  
 033f1cbe55997f39cb331dad525bb7f3  
 cd2406b042aa7f39cb351ca5525bbac4

$SX[K[1]](m)$  = 8d4f93828747a76c49e204adc8473bd1  
 1101dda7470a415b832b77ad5dbc572d  
 111f14950ce8570be4aec9f0e472fd2  
 d9e231ad2c38570be46a14000e47a586

$PSX[K[1]](m)$  = 8d49118311e4d9e44fe2012b1faee26a  
 9304dd7714cd311482ada7ad959fad00  
 87c8475d0c0e2c0e47470abce8473847  
 a73b4157572f57a56cd15b2d0bd20b86

$LPSX[K[1]](m)$  = a3a72a2e0fb5e6f812681222fec037b0  
 db972086a395a387a6084508cae13093  
 aa71d352dcfce288e9a39718a727f6fd  
 4c5da5d0bc10fac3707ccd127fe45475

$K[1] \text{ (xor) } C[1]$  = 92cdb59aaeb185fcc80ec1c1701e230a  
 0caf98039e3e8f03528b56cdc5fe9be9  
 68b90ed1221c36148187c448141b8c00  
 26b39a767c0f1236fe458b1942dd1a12

$S(K[1] \text{ (xor) } C[1])$  = ecd95e282645a83930045858325f5afa  
 2341dc110ad303110ef676d9ac63509b  
 f3a3041b65148f93f5c986f293bb7cfc  
 ef92288ac34df08f63c8f6362cd8f1f0

$PS(K[1] \text{ (xor) } C[1])$  = ec30230ef3f5ef63d90441f6a3c992c8  
 5e58dc76048628f6285811d91bf28a36  
 26320aac6593c32c455fd36314bb4dd8  
 a85a03508f7cf0f139fa119b93fc8ff0

$LPS(K[1] \text{ (xor) } C[1])$  = 18ee8f3176b2e3ea3bd6cb8233694cea  
 349769df88be26bf451cfab6a904a549  
 da22de93a66a66b19c7e6b5eea633511  
 e611d68c8401bfcd0c7d0cc39d4a5eb9

## Iteration 2

K[2] = 18ee8f3176b2ebea3bd6cb8233694cea  
349769df88be26bf451cfab6a904a549  
da22de93a66a66b19c7e6b5eea633511  
e611d68c8401bfcd0c7d0cc39d4a5eb9

LPSX[K[2]]LPSX[K[1]](m) = 9f50697b1d9ce23680db1f4d35629778  
864c55780727aa79eb7bb7d648829cba  
8674afdac5c62ca352d77556145ca7bc  
758679fbel1fbd32313ca8268a4a603f1

## Iteration 3

K[3] = aaa4cf31a265959157aec8ce91e7fd46  
bf27dee21164c5e3940bba1a519e9d1f  
ce0913f1253e7757915000cd674be12c  
c7f68e73ba26fb00fd74af4101805f2d

LPSX[K[3]]...LPSX[K[1]](m) = 4183027975b257e9bc239b75c977ecc5  
2ddad82c091e694243c9143a945b4d85  
3116eae14fd81b14bb47f2c06fd283cb  
6c5e61924edfaf971b78d771858d5310

## Iteration 4

K[4] = 61fe0a65cc177af50235e2afadded326  
a5329a2236747bf8a54228aeca9c4585  
cd801ea9dd743a0d98d01ef0602b0e33  
2067fb5ddd6ac1568200311920839286

LPSX[K[4]]...LPSX[K[1]](m) = 0368c884fcee489207b5b97a133ce39a  
1ebfe5a3ae3cccb3241de1e7ad72857e  
76811d324f01fd7a75e0b669e8a22a4d  
056ce6af3e876453a9c3c47c767e5712

## Iteration 5

K[5] = 9983685f4fd3636f1fd5abb75fbf26a8  
e2934314aa2ecb3ee4693c86c06c7d4e  
169bd540af75e1610a546acd63d960ba  
d595394cc199bf6999a5d5309fe73d5a

LPSX[K[5]]...LPSX[K[1]](m) = c31433ceb8061e46440144e655539765  
12e5a9806ac9a2c771d5932d5f6508c5  
b78e406c4efab98ac5529be0021b4d58  
fa26f01621eb10b43de4c4c47b63f615

## Iteration 6

K[6] = f05772ae2ce7f025156c9a7fbcc6b8fd  
f1e735d613946e32922994e52820ffea  
62615d907eb0551ad170990a86602088  
af98c83c22cdb0e2be297c13c0f7a156

LPSX[K[6]]...LPSX[K[1]](m) = 5d0ae97f252ad04534503fe5f52e9bd0  
7f483ee3b3d206beadc6e736c6e754bb  
713f97ea7339927893eacf2b474a482c  
add9ac2e58f09bcb440cf36c2d14a9b6

## Iteration 7

K[7] = 5ad144c362546e4e46b3e7688829fbb7  
7453e9c3211974330b2b8d0e6be2b5ac  
c89eb6b35167f159b7b005a43e5959a6  
51a9b18cfc8e4098fcf03d9b81cfbb8d

LPSX[K[7]]...LPSX[K[1]](m) = a59aa21e6ad3e330deedb9ab9912205c  
355b1c479fdfd89a7696d7de66fbf7d3  
cec25879f7f1a8cca4c793d5f2888407  
aecb188bda375eae586a8cfd0245c317

## Iteration 8

K[8] = 6a6cec9a1ba20a8db64fa840b934352b  
518c638ed530122a83332fe0b8efdac9  
018287e5a9f509c78d6c746adcd5426f  
b0a0ad5790dfb73fc1f191a539016daa

LPSX[K[8]]...LPSX[K[1]](m) = 9903145a39d5a8c83d28f70fa1fbd88f  
31b82dc7cfe17b54b50e276cb2c4ac68  
2b4434163f214cf7ce6164a75731bcea  
5819e6a6a6fea99da9222951d2a28e01

## Iteration 9

K[9] = 99217036737aa9b38a8d6643f705bd51  
f351531f948f0fc5e35fa35fee9dd8bd  
bb4c9d580a224e9cd82e0e2069fc49ed  
367d5f94374435382b8fb6a8f5dd0409

LPSX[K[9]]...LPSX[K[1]](m) = 330e6cb1d04961826aa263f2328f15b4  
f3370175a6a9fd6505b286efed2d8505  
f71823337ef71513e57a700eb1672a68  
5578e45dad298ee2223d4cb3fda8262f

Iteration 10

K[10] = 906763c0fc89fa1ae69288d8ec9e9dda  
9a7630e8bfd6c3fed703c35d2e62aeaf  
f0b35d80a7317a7f76f83022f2526791  
ca8fdf678fcb337bd74fe5393ccb05d2

LPSX[K[10]]...LPSX[K[1]](m) = ad347608443ab9c9bbb64f633a5749ab  
85c45d4174bfd78f6bc79fc4f4ce9ad1  
dd71cb2195b1cfab8dcaaf6f3a65c8bb  
0079847a0800e4427d3a0a815f40a644

Iteration 11

K[11] = 88ce996c63618e6404a5c8e03ee43385  
4e2ae3eee68991bbbff3c29d38dadb6e  
d6a1dae9a6dc6ddf52ce34af272f96d3  
159c8c624c3fe6e13d695c0bfc89add5

LPSX[K[11]]...LPSX[K[1]](m) = a065c55e2168c31576a756c7ecc1a912  
9cd3d207f8f43073076c30e111fd5f11  
9095ca396e9fb78a2bf4781c44e845e4  
47b8fc75b788284aae27582212ec23ee

Iteration 12

K[12] = 3e0a281ea9bd46063eec550100576f3a  
506aa168cf82915776b978fcca32f38  
b55f30c79982ca45628e8365d8798477  
e75a49c68199112a1d7b5a0f7655f2db

LPSX[K[12]]...LPSX[K[1]](m) = 2a6549f7a5cd2eb4a271a7c71762c868  
3e7a3a906985d60f8fc86f64e35908b2  
9f83b1fe3c704f3c116bdf660704f3b  
9c8a1d0531baaffaa3940ae9090a33ab

Iteration 13

K[13] = f0b273409eb31aeb432fbae18672122  
62c848422b6a92f93f6cbab54ed18b83  
14b21cffc51e3fa319ff433e76ef6adb  
0ef9f5e03c907fa1fcf9eca06500bf03

X[K[13]]...LPSX[K[1]](m) = dad73ab73b7e345f46435c690f05e94a  
5cb272d242ef44f6b0a4d5d1ad888331  
8b31ad01f96e709f08949cd8169f25e0  
9273e8e50d2ad05b5f6de6496c0a8ca8



The result of the transformation  $g_0(h, N)$  is

```
h = aee3bd55ea6f387bcf28c6dcdbbfb3d
    dacc67dcc13dbd8d548c6bf808111d4b
    75b8e74d2afae960835ae6a5f0357555
    9c9fd839783ffcd5cf99bd61566b4818
```

The result of the transformation  $g_0(h, \text{EPSILON})$  is

```
h = 508f7e553c06501d749a66fc28c6cac0
    b005746d97537fa85d9e40904efed29d
    c345e53d7f84875d5068e4eb743f0793
    d673f09741f9578471fb2598cb35c230
```

The hash code of the message M2 is the value

```
H(M2) = 508f7e553c06501d749a66fc28c6cac0
        b005746d97537fa85d9e40904efed29d
```

## 11. Security Considerations

This entire document is about security considerations.

## 12. References

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