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K. Sriram
D. Montgomery
US NIST
D. McPherson
E. Osterweil
Verisign, Inc.
B. Dickson
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Problem Definition and Classification of BGP Route Leaks

Abstract

A systemic vulnerability of the Border Gateway Protocol routing system, known as "route leaks", has received significant attention in recent years. Frequent incidents that result in significant disruptions to Internet routing are labeled route leaks, but to date a common definition of the term has been lacking. This document provides a working definition of route leaks while keeping in mind the real occurrences that have received significant attention. Further, this document attempts to enumerate (though not exhaustively) different types of route leaks based on observed events on the Internet. The aim is to provide a taxonomy that covers several forms of route leaks that have been observed and are of concern to the Internet user community as well as the network operator community.

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1. Introduction

Frequent incidents [[Huston2012](#)] [[Cowie2013](#)] [[Toonk2015-A](#)] [[Toonk2015-B](#)] [[Cowie2010](#)] [[Madory](#)] [[Zmijewski](#)] [[Paseka](#)] [[LRL](#)] [[Khare](#)] that result in significant disruptions to Internet routing are commonly called "route leaks". Examination of the details of some of these incidents reveals that they vary in their form and technical details. In order to pursue solutions to "the route-leak problem" it is important to first provide a clear, technical definition of the problem and enumerate its most common forms. [Section 2](#) provides a working definition of route leaks, keeping in view many recent incidents that have received significant attention. [Section 3](#) attempts to enumerate (though not exhaustively) different types of route leaks based on observed events on the Internet. Further, [Section 3](#) provides a taxonomy that covers several forms of route leaks that have been observed and are of concern to the Internet user community as well as the network operator community. This document builds on and extends earlier work in the IETF [[ROUTE-LEAK-DEF](#)] [[ROUTE-LEAK-REQ](#)].

2. Working Definition of Route Leaks

A proposed working definition of "route leak" is as follows:

A route leak is the propagation of routing announcement(s) beyond their intended scope. That is, an announcement from an Autonomous System (AS) of a learned BGP route to another AS is in violation of the intended policies of the receiver, the sender, and/or one of the ASes along the preceding AS path. The intended scope is usually defined by a set of local redistribution/filtering policies distributed among the ASes involved. Often, these intended policies are defined in terms of the pair-wise peering business relationship between ASes (e.g., customer, transit provider, peer). For literature related to AS relationships and routing policies, see [[Gao](#)], [[Luckie](#)], and [[Gill](#)]. For measurements of valley-free violations in Internet routing, see [[Anwar](#)], [[Giotsas](#)], and [[Wijchers](#)].

The result of a route leak can be redirection of traffic through an unintended path that may enable eavesdropping or traffic analysis and may or may not result in an overload or black hole. Route leaks can be accidental or malicious but most often arise from accidental misconfigurations.

The above definition is not intended to be all encompassing. Our aim here is to have a working definition that fits enough observed incidents so that the IETF community has a basis for developing solutions for route-leak detection and mitigation.

3. Classification of Route Leaks Based on Documented Events

As illustrated in Figure 1, a common form of route leak occurs when a multihomed customer AS (such as AS3 in Figure 1) learns a prefix update from one transit provider (ISP1) and leaks the update to another transit provider (ISP2) in violation of intended routing policies, and further, the second transit provider does not detect the leak and propagates the leaked update to its customers, peers, and transit ISPs.

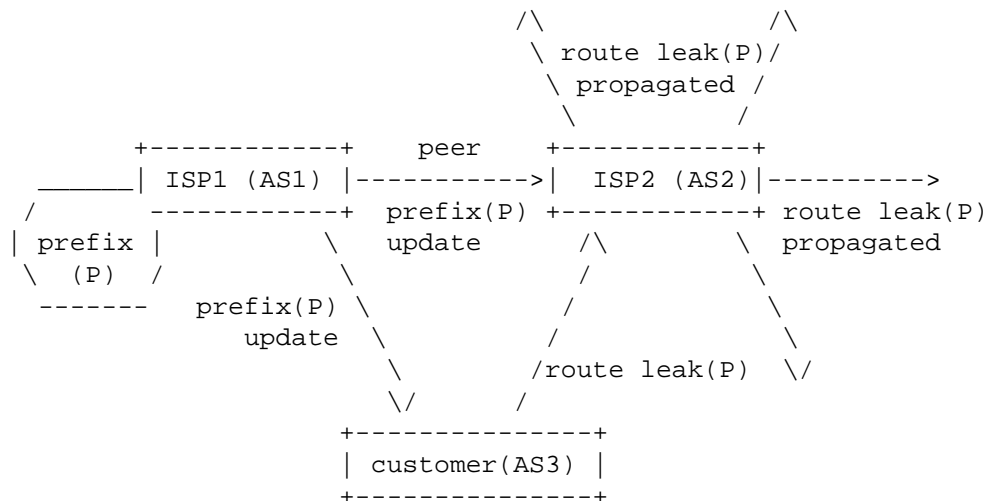


Figure 1: Basic Notion of a Route Leak

This document proposes the following taxonomy to cover several types of observed route leaks while acknowledging that the list is not meant to be exhaustive. In what follows, the AS that announces a route that is in violation of the intended policies is referred to as the "offending AS".

3.1. Type 1: Hairpin Turn with Full Prefix

Description: A multihomed AS learns a route from one upstream ISP and simply propagates it to another upstream ISP (the turn essentially resembling a hairpin). Neither the prefix nor the AS path in the update is altered. This is similar to a straightforward path-poisoning attack [Kapela-Pilosov], but with full prefix. It should be noted that leaks of this type are often accidental (i.e., not malicious). The update basically makes a hairpin turn at the offending AS's multihomed AS. The leak often succeeds (i.e., the leaked update is accepted and propagated) because the second ISP prefers customer announcement over peer announcement of the same prefix. Data packets would reach the legitimate destination, albeit

via the offending AS, unless they are dropped at the offending AS due to its inability to handle resulting large volumes of traffic.

- o Example incidents: Examples of Type 1 route-leak incidents are (1) the Dodo-Telstra incident in March 2012 [[Huston2012](#)], (2) the VolumeDrive-Atrato incident in September 2014 [[Madory](#)], and (3) the massive Telekom Malaysia route leak of about 179,000 prefixes, which in turn Level3 accepted and propagated [[Toonk2015-B](#)].

3.2. Type 2: Lateral ISP-ISP-ISP Leak

Description: The term "lateral" here is synonymous with "non-transit" or "peer-to-peer". This type of route leak typically occurs when, for example, three sequential ISP peers (e.g., ISP-A, ISP-B, and ISP-C) are involved, and ISP-B receives a route from ISP-A and in turn leaks it to ISP-C. The typical routing policy between laterally (i.e., non-transit) peering ISPs is that they should only propagate to each other their respective customer prefixes.

- o Example incidents: In [[Mauch-nanog](#)] and [[Mauch](#)], route leaks of this type are reported by monitoring updates in the global BGP system and finding three or more very large ISPs' Autonomous System Numbers (ASNs) in a sequence in a BGP update's AS path. [[Mauch](#)] observes that its detection algorithm detects for these anomalies and potentially route leaks because very large ISPs do not, in general, buy transit services from each other. However, it also notes that there are exceptions when one very large ISP does indeed buy transit from another very large ISP, and accordingly, exceptions are made in its detection algorithm for known cases.

3.3. Type 3: Leak of Transit-Provider Prefixes to Peer

Description: This type of route leak occurs when an offending AS leaks routes learned from its transit provider to a lateral (i.e., non-transit) peer.

- o Example incidents: The incidents reported in [[Mauch](#)] include Type 3 leaks.

3.4. Type 4: Leak of Peer Prefixes to Transit Provider

Description: This type of route leak occurs when an offending AS leaks routes learned from a lateral (i.e., non-transit) peer to its (the AS's) own transit provider. These leaked routes typically originate from the customer cone of the lateral peer.

- o Example incidents: Examples of Type 4 route-leak incidents are (1) the Axcclx-Hibernia route leak of Amazon Web Services (AWS) prefixes causing disruption of AWS and a variety of services that run on AWS [[Kephart](#)], (2) the Hathway-Airtel route leak of 336 Google prefixes causing widespread interruption of Google services in Europe and Asia [[Toonk2015-A](#)], (3) the Moratel-PCCW route leak of Google prefixes causing Google's services to go offline [[Paseka](#)], and (4) some of the example incidents cited for Type 1 route leaks above are also inclusive of Type 4 route leaks. For instance, in the Dodo-Telstra incident [[Huston2012](#)], the leaked routes from Dodo to Telstra included routes that Dodo learned from its transit providers as well as lateral peers.

3.5. Type 5: Prefix Re-origination with Data Path to Legitimate Origin

Description: A multihomed AS learns a route from one upstream ISP and announces the prefix to another upstream ISP as if it is being originated by it (i.e., strips the received AS path and re-originate the prefix). This can be called re-origination or mis-origination. However, somehow a reverse path to the legitimate origination AS may be present and data packets reach the legitimate destination albeit via the offending AS. (Note: The presence of a reverse path here is not attributable to the use of a path-poisoning trick by the offending AS.) But sometimes the reverse path may not be present, and data packets destined for the leaked prefix may be simply discarded at the offending AS.

- o Example incidents: Examples of Type 5 route leak include (1) the China Telecom incident in April 2010 [[Hiran](#)] [[Cowie2010](#)] [[Labovitz](#)], (2) the Belarusian GlobalOneBel route-leak incidents in February-March 2013 and May 2013 [[Cowie2013](#)], (3) the Icelandic Opin Kerfi-Simmin route-leak incidents in July-August 2013 [[Cowie2013](#)], and (4) the Indosat route-leak incident in April 2014 [[Zmijewski](#)]. The reverse paths (i.e., data paths from the offending AS to the legitimate destinations) were present in incidents #1, #2, and #3 cited above, but not in incident #4. In incident #4, the misrouted data packets were dropped at Indosat's AS.

3.6. Type 6: Accidental Leak of Internal Prefixes and More-Specific Prefixes

Description: An offending AS simply leaks its internal prefixes to one or more of its transit-provider ASes and/or ISP peers. The leaked internal prefixes are often more-specific prefixes subsumed by an already announced, less-specific prefix. The more-specific prefixes were not intended to be routed in External BGP (eBGP). Further, the AS receiving those leaks fails to filter them.

Typically, these leaked announcements are due to some transient failures within the AS; they are short-lived and typically withdrawn quickly following the announcements. However, these more-specific prefixes may momentarily cause the routes to be preferred over other aggregate (i.e., less specific) route announcements, thus redirecting traffic from its normal best path.

- o Example incidents: Leaks of internal routes occur frequently (e.g., multiple times in a week), and the number of prefixes leaked range from hundreds to thousands per incident. One highly conspicuous and widely disruptive leak of internal routes happened in August 2014 when AS701 and AS705 leaked about 22,000 more-specific prefixes of already-announced aggregates [Huston2014] [Toonk2014].

4. Additional Comments about the Classification

It is worth noting that Types 1 through 4 are similar in that a route is leaked in violation of policy in each case, but what varies is the context of the leaked-route source AS and destination AS roles.

A Type 5 route leak (i.e., prefix mis-origination with data path to legitimate origin) can also happen in conjunction with the AS relationship contexts in Types 2, 3, and 4. While these possibilities are acknowledged, simply enumerating more types to consider all such special cases does not add value as far as solution development for route leaks is concerned. Hence, the special cases mentioned here are not included in enumerating route-leak types.

5. Security Considerations

No security considerations apply since this is a problem definition document.

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Authors' Addresses

Kotikalapudi Sriram
US NIST

Email: ksriram@nist.gov

Doug Montgomery
US NIST

Email: dougmon@nist.gov

Danny McPherson
Verisign, Inc.

Email: dmcpherson@verisign.com

Eric Osterweil
Verisign, Inc.

Email: eosterweil@verisign.com

Brian Dickson

Email: brian.peter.dickson@gmail.com