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- In all cases, evaluating the conditional expression results in executing one of two other pieces of code identified by the symbolic labels L_{TRUE} and L_{FALSE} .

Conditionals: Short-Circuit Evaluation

- Evaluating $E = (E1 \ \& \ E2)$.
 - Evaluate E1.
 - If E1 evaluates to false, go to L_{FALSE} ; otherwise,
 - Evaluate E2.
 - If E2 evaluates to true, go to L_{TRUE} ; otherwise, go to L_{FALSE} .
- Evaluating $E = (E1 \ | \ E2)$.
 - Evaluate E1.
 - If E1 evaluates to true, go to L_{TRUE} ; otherwise,
 - Evaluate E2.
 - If E2 evaluates to true, go to L_{TRUE} ; otherwise, go to L_{FALSE} .
- Evaluating $E = (! \ E1)$.
 - Evaluate E1.
 - If E1 evaluates to true, go to L_{FALSE} ; otherwise, go to L_{TRUE} .
- The “Evaluate E1” and “Evaluate E2” steps may themselves be recursive.

Conditionals: Short-Circuit Code Generation

- $codegen(E, L_{TRUE}, L_{FALSE})$, where $E = (E1 \ \& \ E2)$.
 - Let $L_{new} = newlabel()$ and $L_{end} = newlabel()$.
 - $C1 = codegen(E1, L_{new}, L_{FALSE})$.
 - $C2 = codegen(E2, L_{TRUE}, L_{FALSE})$.
 - **return** $C1 + "L_{new}:\backslash n" + C2 + "L_{TRUE}:\backslash n" + \text{PUSHIMM } 1\backslash n + \text{JUMP } L_{end}\backslash n + "L_{FALSE}:\backslash n" + \text{PUSHIMM } 0\backslash n + "L_{end}:\backslash n"$.

Conditionals: Short-Circuit Code Generation

- *codegen*(*E*, *L_{TRUE}*, *L_{FALSE}*), where *E* = (*E1* & *E2*).
 - Let *L_{new}* = *newlabel*() and *L_{end}* = *newlabel*().
 - *C1* = *codegen*(*E1*, *L_{new}*, *L_{FALSE}*).
 - *C2* = *codegen*(*E1*, *L_{TRUE}*, *L_{FALSE}*).
 - **return** *C1* + "*L_{new}*:
" + *C2* + "*L_{TRUE}*:
" + "PUSHIMM 1
" + "JUMP
L_{end}
" + "*L_{FALSE}*:
" + "PUSHIMM 0
" + "*L_{end}*:
".
- *codegen*(*E*, *L_{TRUE}*, *L_{FALSE}*), where *E* = (*E1* | *E2*).
 - *C1* = *codegen*(*E1*, *L_{TRUE}*, *L_{new}*), where *L_{new}* = *newlabel*().
 - *C2* = *codegen*(*E1*, *L_{TRUE}*, *L_{FALSE}*).
 - **return** *C1* + "*L_{new}*:
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 - Let *L_{new}* = *newlabel*() and *L_{end}* = *newlabel*().
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 - **return** *C1* + "*L_{new}*:*\n*" + *C2* + "*L_{TRUE}*:*\n*" + "PUSHIMM 1*\n*" + "JUMP *L_{end}**\n*" + "*L_{FALSE}*:*\n*" + "PUSHIMM 0*\n*" + "*L_{end}*:*\n*".
- *codegen*(*E*, *L_{TRUE}*, *L_{FALSE}*), where *E* = (*E1* | *E2*).
 - *C1* = *codegen*(*E1*, *L_{TRUE}*, *L_{new}*), where *L_{new}* = *newlabel*().
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 - **return** *C1* + "*L_{new}*:*\n*" + *C2* + "*L_{TRUE}*:*\n*" + "PUSHIMM 1*\n*" + "JUMP *L_{end}**\n*" + "*L_{FALSE}*:*\n*" + "PUSHIMM 0*\n*" + "*L_{end}*:*\n*".
- *codegen*(*E*, *L_{TRUE}*, *L_{FALSE}*), where *E* = (! *E1*).
 - **return** *codegen*(*E1*, *L_{FALSE}*, *L_{TRUE}*).

Example of Short-Circuit Code Generation

```
codegen(( (a<b) | ((c<d) & (e<f)) ), Ltrue, Lfalse)
```