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- The code aggregates malicious computing by linking together short code snippets already present in the program's address space.
 - Each snippet ends with a ret instruction, which allows an attacker controlling the stack to chain them together.
 - Because the executed code is stored in memory marked executable, page protection mechanisms won't help.

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 - Because the executed code is stored in memory marked executable, page protection mechanisms won't help.
- Still needs a buffer overflow or similar mechanism to subvert the control flow of the program.

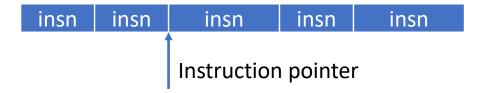
Program Layout

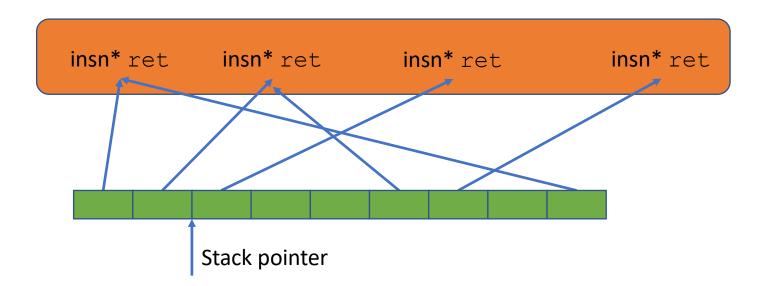
- An ordinary program (OP) is a sequence of machine instructions (MIs) laid out in the text segment of the program.
 - Each MI is a byte pattern that, when interpreted by the processor, induces a small amount of change in machine state.
 - The instruction pointer %rip governs the MI that is fetched and executed next.

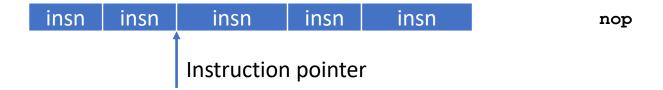
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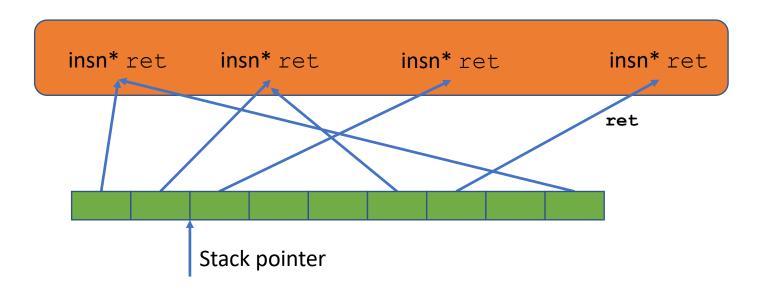
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 - Each MI is a byte pattern that, when interpreted by the processor, induces a small amount of change in machine state.
 - The instruction pointer %rip governs the MI that is fetched and executed next.
- A return-oriented program (RP) is a sequence of returnoriented instructions (RIs) laid out in the *stack* segment of the exploited program.
 - Each RI is a sequence of words on the stack pointing to a sequence of MIs ending in a ret, somewhere in the address space of the exploited program.
 - The stack pointer governs what MI sequence is to be fetched next in the following way. The execution of a ret instruction has two effects: first, the word to which <code>%rsp</code> points is read and used as the new value for <code>%rip</code>; second, <code>%rsp</code> is incremented by 8 bytes to point to the next word on the stack.

insn	insn	insn	insn	insn
<u> </u>				
		Instruction pointer		



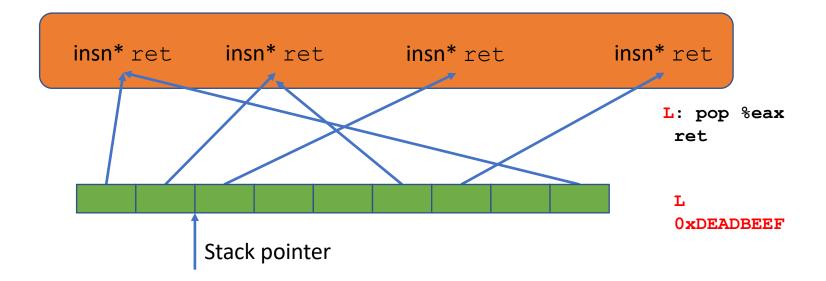








mov \$0xDEADBEEF, %eax



Gadgets

- A gadget is an arrangement of words on the stack, including one or more instruction sequence pointers and associated immediate values, that encodes a logical unit of work in the RP.
 - Gadgets act like a return-oriented instruction set and are the natural target of a return-oriented compiler's assembler.

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 - Gadgets act like a return-oriented instruction set and are the natural target of a return-oriented compiler's assembler.
- Example: A memory-load gadget.
 - MI sequences
 - S_1 : pop %eax; ret
 - S_2 : mov (%eax), %ebx; ret
 - On stack

[Pointer to S_1 , address of word to load, pointer to S_2]

Gadget Execution

- Correct execution of a gadget
 - Precondition: %rsp points to the first word in the gadget and the processor executes a ret instruction.
 - Postcondition: when the ret instruction in its last instruction sequence is executed, %rsp points to the next gadget to be executed.

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Exploitation

- Place the payload containing these gadgets in the memory of the program to be exploited, and redirect the stack pointer so that it points to the first gadget.
- The easiest way to accomplish these tasks is by means of a buffer overflow on the stack; the gadgets are placed on the overflowed stack so that the first has overwritten the saved instruction pointer of some function. When that function tries to return, the returnoriented program is executed instead.

Finding Useful Instruction Sequences

- Every instruction sequence ending in a ret instruction (represented by the byte $0 \times c3$ on x86) is potentially useful.
- The suffix of any useful instruction sequence is also useful.
- The frequency of occurrence of some sequence is not relevant; only its existence is.

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- The frequency of occurrence of some sequence is not relevant; only its existence is.
- On variable-length instruction sequences, there can also be unintended instruction sequences.

```
f7 c7 07 00 00 00 test $0x00000007, %edi
0f 95 45 c3 setnzb -61(%ebp)

c7 07 00 00 00 0f movl $0x0f000000, (%edi)
95 xchg %ebp, %eax
inc %ebp
c3
```