

## How I solve the problem

Assume I have 4 processors, and  $N = 16$ . Initially they generate list of numbers individually.

P0: 2 3 5 7  
P1: 4 5 1 2  
P2: 9 3 2 1  
P3: 7 7 6 6

We randomly select the pivot from each processor

P1: 2, P2: 5, P3 9, P4: 6

Use `MPI_Gather()` and processor 0 gathers all of the pivots. Processor 0 sort it and broad the medium 6 using `MPI_Bcast`. Now every processor knows they should part the data with pivot 6.

P0: 2 3 5 / 7  
P1: 4 5 1 2 /  
P2: 3 2 1 / 9  
P3: 6 6 / 7 7

The data in each processors are divided into two part,  $<$  and  $\geq$  part.  
Then we divide the communicators( processor) into two groups. (P0,P1) and (P2,P3).  
(P2,P3) send its  $<$  part to (P0,P1) and (P0,P1) send its  $\geq$  part to (P2,P3). Specifically  
P0 - P2, P1 - P3.

This part using `MPI_Send` and `MPI_Recv`.

Now

P0: 2 3 5 3 2 1  
P1: 4 5 1 2 6 6  
  
P2: 9 7  
P3: 7 7

Then recursively do the previous procedure in each group, until each group only have one communicator. Then naively quick sort it.

After quick sort. Every processor hold part of the array in order. And  
 $P0 \leq P1 \leq P2 \leq P3$ . So I finished the problem. Then use `MPI_Send`& `MPI_Recv` to

gather all of the data in processor with world rank 0. And Processor 0 use MPI\_Scatter to make sure each processor has  $N/P$  elements.

## Excution Time

	$N = 10^6$	$N = 10^7$	$N = 10^8$
$P = 1$	0.2813s	3.0506s	37.0299s
$P = 2$	0.3008s	2.1368s	28.4473s
$P = 4$	0.2804s	2.4349s	18.5802s
$P = 8$	0.0839s	1.0315s	15.6301s
$P = 16$	0.1235s	1.2213s	8.6602s