#### Linux Conference

October 26- 28, 2011 Prague . Czech Republic



## The Extended FAT file system

Differentiating with FAT32 file system

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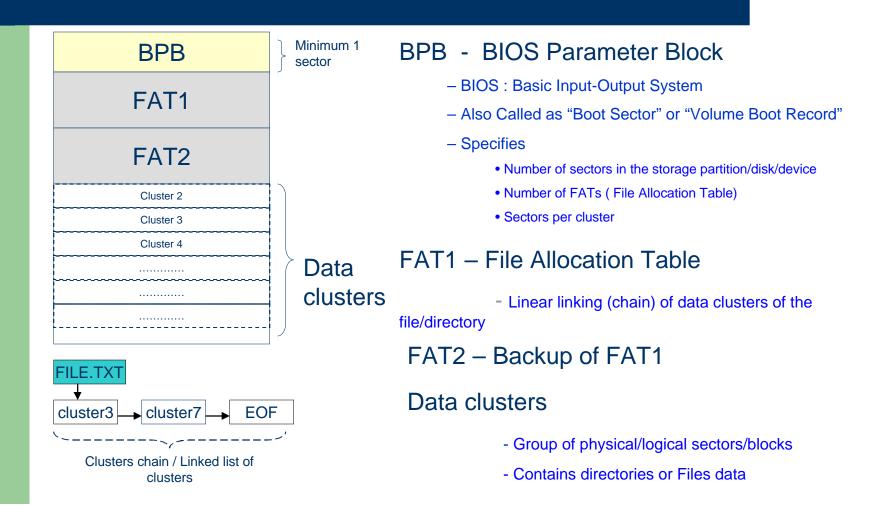


### Agenda

- Ø FAT file system
- Ø Need for ExFAT file system
- Ø ExFAT file system organization
- Ø Boot Sector of FAT and ExFAT
- Ø Snapshot of Boot Sector of ExFAT
- Ø Directory Entries of FAT
- Ø Directory Entries of ExFAT
- Ø Clusters Heap
- Ø Up-Case Table
- Ø Snapshot of Root directory
- Ø Performance Benchmarking in Linux
- Ø Performance Benchmarking in Windows
- Ø References



### **FAT File system**



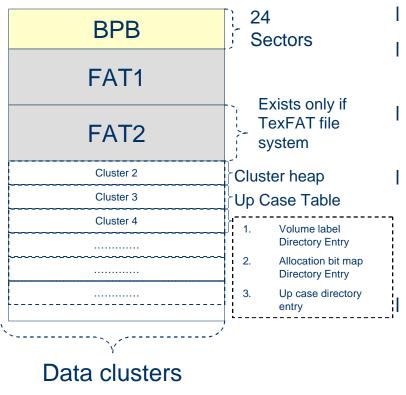


### **Need For ExFAT file system**

- FAT File system limited to support only 32GB.
  - FAT does not support Higher size SDXC cards.
- **I** NTFS
  - Security features
    - Optional for Removable storage devices
  - Meta data overhead for file/directory
  - Write caching mechanism for performance optimizations causes the data corruption in when removable storage device is unplugged.



### **ExFAT file system Organization**



- 12 Sectors for Boot Area
- Another 12 sectors for backup of Boot Area
  - FAT2 exists if the TexFAT ( Transactional FAT is used)
    - First few Data clusters starting from cluster 2 contains
      - Cluster allocation bit map table
      - Up case table

#### **Next Cluster Contains**

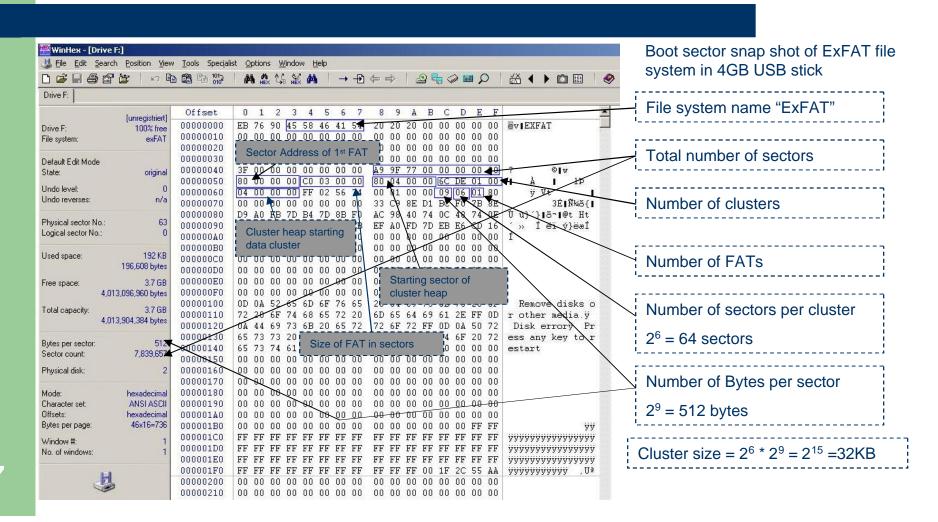
- Cluster allocation Bit map Directory Entry
- Up-case Table Directory Entry
- Volume label Directory Entry



#### **Key Elements of Boot Sector of ExFAT and FAT32**

Field name		E	xFAT Boot Sector	FAT32 Boot Sector			
	Offset ( Byte)	Size ( Byte)	Description	Offset (Byte)	Size ( Byte)	Description	
File system name	3	8	"ExFAT"	"MSWIN4	.1" or "MSDC	DS5.0"	
				82	8	"FAT32 "	
Volume length	72	8	Total number of Sectors	32	4	Total number of Sectors	
FAT offset	80	4	Sector address of 1st FAT				
FAT length	84	4	Size of FAT in sectors	36	4	Size of 1st FAT in sectors	
Cluster offset	88	4	Starting sector of cluster heap				
Cluster count	92	4	Number of clusters				
Root directory	96	4	First cluster of root directory	44	4	First cluster of root directory	
Volume flags	106	2	Bit 0 – Active FAT  0 – 1st , 1 – 2nd  Bit 1 – Volume Dirty  0 – Clean, 1- dirty  Bits 2 & 3 – Media failure  0 – No failures, 1 – failures reported				
Bytes per Sector	108	1	This is power of 2; Minimal value is 9; 29 =512 Bytes and maximum 212 =4096 Bytes	11	2	Count of bytes per sector; This can have following values 512,1024,2048 or 4096.	
Sector per cluster	109	1	This is power of 2; Minimal value is 1; 2° =1 sector (512 Bytes) and maximum 32 MB cluster size in bytes	13	1	Number of sectors per cluster; allowed values are 1,2,4,8,32,64 and 128. Maximum of 128 sectors per cluster	
Number of FATs	110	1	Either 1 or 2; if TexFAT is supported then it will be 2	16	1	Usually 2	

#### **Boot Sector of ExFAT**



# **Default Cluster size of ExFAT, FAT and NTFS**

Disk size	FAT16	FAT32	NTFS	ExFAT
4 MB – 16MB	1KB	512 Bytes	4KB	4KB
16 MB-64MB	2KB			
64 MB-128 MB				
128 MB-256 MB	4KB			
256MB- 512MB	8KB	4KB		32KB
512MB- 1GB	16KB			
1GB- 2GB	32KB			
2GB – 8GB	Not supported			
8GB-16GB		8KB		
16GB-32GB		16KB		
32GB- 127GB		32KB		128KB
127GB-2TB		Not supported;		
2TB-16TB		*Assuming Sector size is 512		
16TB – 32TB		bytes, then beyond 127GB is not supported by FAT32, as	8KB	
32TB-64TB		per FAT spec, the FAT entry size is 32 bits but only 28 bits	16KB	
64TB – 128TB		are used for cluster entry.	32KB	
128TB – 256TB			64KB	
>256TB			Custom specific	Custom specific

Note: \*

KB – Kilo Bytes

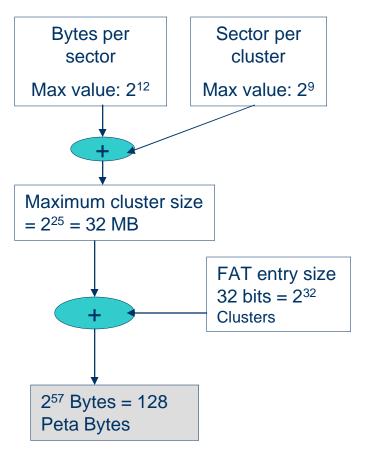
MB – Mega Bytes

GB - Giga Bytes

TB – Tera Bytes



### **ExFAT Maximum Volume Support**



- The "Volume length" field of Boot sector is 8 bytes; so the maximum value is 2<sup>64</sup> Sectors
- The Maximum value of "Bytes per sector" field of boot sector is  $2^{12} = 4096$  bytes
- So, the Maximum theoretical value support is  $2^{(64+12)} = 2^{76}$  which is 64 Zetta Bytes (ZB)
- The FAT entry size is 32, Maximum cluster number is 2<sup>32</sup>
- So, To support 64 ZB, it requires at least  $2^{(32+12)}$  16 Terra Bytes (TB) cluster size.
  - But, Microsoft has set the limitation that cluster size should be of maximum of 32 MB
    - Hence "Bytes per sector" + "Sectors per cluster" of boot sector can not exceed 25. hence the maximum value will be  $2^{25}$  and the maximum cluster are  $2^{32}$ , so it yields maximum value of  $2^{(32+25)} = 2^{57}$  Bytes = 128 Peta Bytes (PB)
- Practically, Maximum Supported Size is 128 PB



## **Directory Entries of FAT**

32 Bytes Directory Entry of FAT32					
Name	Offset (byte)	Size (bytes)	Description		
DIR_Name	0	11	Short name.		
DIR_Attr	11	1	File attributes		
DIR_NTRes	12	1	Reserved		
DIR_CrtTimeTenth	13	1	Millisecond stamp at file creation time		
DIR_CrtTime	14	2	Time file was created.		
DIR_CrtDate	16	2	Date file was created.		
DIR_LstAccDate	18	2	Last access date.		
DIR_FstClusHI	20	2	High word of this entry's first cluster number		
DIR_WrtTime	22	2	Time of last write.		
DIR_WrtDate	24	2	Date of last write.		
DIR_FstClusLO	26	2	Low word of this entry's first cluster number.		
DIR_FileSize	28	4	file size in bytes.		

F	FAT Long Directory Entry Structure						
Name	Offset (byte)	Size (bytes)	Description				
LDIR_Ord	0	1	Masked with 0x40 (LAST_LONG_ENTRY), this indicates the entry is the last long dir entry in a set of long dir entries. All valid sets of long dir entries must begin with an entry having this mask.				
LDIR_Name1	1	10	Characters 1-5 of the long-name sub- component in this dir entry.				
LDIR_Attr	11	1	Attributes - must be ATTR_LONG_NAME				
LDIR_Type	12	1	Zero				
LDIR_Chksum	13	1	Checksum of name in the short dir entry at the end of the long dir set.				
LDIR_Name2	14	12	Characters 6-11 of the long-name sub- component in this dir entry.				
LDIR_FstClusLO	26	2	Must be ZERO.				
LDIR_Name3	28	4	Characters 12-13 of the long-name sub- component in this dir entry.				



### **Directory Entries of ExFAT**

- Following are the Directory Entry Types defined
  - Critical Primary
    - Volume Label Directory Entry
    - Allocation Bitmap Directory Entry
    - Up-case Table Directory Entry
    - File/Folder directory Entry
  - Benign primary
    - I Volume GUID
    - TexFAT Padding
    - Windows CE Access Control Table
  - Critical Secondary
    - Stream Extension Directory Entry
    - File name Extension Directory Entry
- All Directory Entries are of 32 Bytes



### **Volume Label Directory Entry**

Name	Offset (byte)	Size (bytes)	Description
Entry Type	0	1	0x83 – Indicates that the Volume label exists 0x03- Indicates that the volume label is deleted
Character count	1	1	Number of characters in label
Volume label	2	22	Volume label in Unicode;

8

24

Maximum of 11 characters of

16 bit Unicode string

Volume Label Directory Entry of ExFAT

Reserved

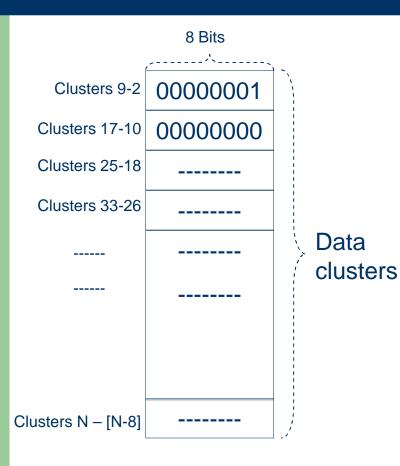


### Allocation Bit map (Cluster Heap) Directory Entry

#### Cluster Heap Directory Entry of ExFAT

Name	Offset (byte)	Size (bytes)	Description
Entry Type	0	1	0x81
Bit Map Flags	1	1	Bit 0 0 - 1 <sup>st</sup> cluster Heap 1 - 2 <sup>nd</sup> cluster Heap
Reserved	2	18	
First cluster	20	4	First data cluster number; Usually it will be 2
Data length	24	8	Size of cluster heap in bytes It is calculated as: Ceil (Cluster count / 8)

### **Cluster Heap**



- Generally the cluster 2 contains the cluster heap
- if cluster heap size is more than cluster size; then the index 2 of File Allocation Table (FAT) contains the next cluster number which contains the cluster heap.
- There can exist multiple clusters ( cluster chain in FAT) as cluster heap
- Every Bit in the Cluster heap contains the status of the data cluster; 0 Free; 1- Allocated
- Byte 0 contains allocation status of cluster 2 to cluster 9
- Byte 1 contains the allocation status of cluster 10 to cluster 17
- Byte 2 contains the allocation status of cluster 18 to cluster 25 and so on

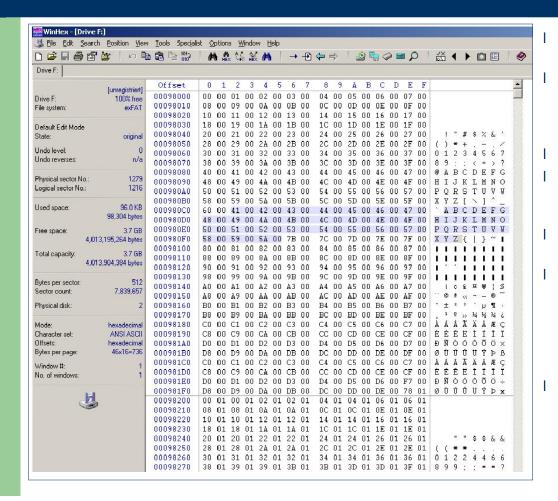


## **Up-Case Table Directory Entry**

Up Case Table Directory Entry of ExFAT					
Name	Offset (byte)	Size (bytes)	Description		
Entry Type	0	1	0x82		
Reserved	1	3			
Table checksum	4	4	Check sum of the Up-case table		
Reserved	8	12			
First cluster	20	4	First data cluster		
Data length	24	8	Up case table size in bytes		



### **Up-Case Table**



The Up case Table is series of Unicode character mappings.

One of the uses of the Up case table is "lower case" to "Upper case" conversion when file/directory searching is performed.

- Up case table is always Read-only.
- If cluster heap is only one cluster and it is cluster 2; then the Up case table starts from cluster 3.
- By default, Up-case table starts immediately after the cluster heap.
- If Up case table size is more than cluster size; File Allocation Table (FAT) contains the cluster chain specifying the data clusters in which content of Up case table is available.

#### Example:

- WinHex tool Snap shot of first few bytes of Up-case Table
- Lower case letter "a ( 0x:61)" converted to Upper case "A (0x41)" by having value 0x0041 at the 16 bit index 0x61.



## **File Directory Entry**

File/Folder Directory Entry of ExFAT					
Name	Offset (byte)	Size (bytes)	Description		
Entry Type	0	1	0x85		
Secondary count	1	1	Number of Secondary directory entries; Minimum value is 2 and maximum value is 18. These entries will be "Stream Extension directory entry" and "File name extension directory entries"; In general, A file/directory contains 1 File directory Entry, 1 stream extension directory and maximum of 17 file name extension directory entries.		
Set checksum	2	2	Checksum is calculated on File directory entry (this entry) + All secondary entries		
File Attributes	4	2	Attributes of the file/directory		
Reserved	6	2			
Creation time	8	4	Created Date and Time in DOS time stamp format		
Last Modified	12	4	Last modified Date and Time in DOS time stamp format		
Last Accessed	16	4	Last accessed Date and Time in DOS time stamp format		
Create 10ms Increment	20	1	10ms increments between 0 – 199		
Last modified 10ms Increment	21	1	10ms increments between 0 – 199		
Last accessed 10ms Increment	22	1	10ms increments between 0 – 199		
Reserved	23	9			



#### File Attributes Field structure

File Attributes field of File Directory Entry of ExFAT						
Name	Offset (byte)	Size (bytes)	Description			
Read only	0	1	0 – Write permission 1 – Read only			
Hidden	1	1	0 – Visible 1 – Hidden			
System	2	1	1 – system			
Reserved	3	1	In FAT32, this bit is used to indicate that the file/directory entry is volume label			
Directory	4	1	0 – File 1 – Directory/Folder			
Archive	5	1	1- Archive			
Reserved	6	10				

- The Read-only, Hidden, System, Directory and Archive bit fields are file attribute are same as FAT32 File attributes structure.
- Volume Id is reserved in file attribute of ExFAT, because there is exist a dedicated directory entry to represent the volume name of the disk/partition



## File/Directory Timestamp Structure

	Timestamp field of File Directory Entry of ExFAT						
Name	Offset (byte)	Size (bytes)	Description				
Double Seconds	0	5	Represents the seconds in two-second multiples.  Valid range of values is:  0 – which represents 0 seconds  29 – represents 58 seconds				
Minute	5	6	Minutes; Valid range 0 -59				
Hour	11	5	Hours; Valid range 0 -23				
Day	16	5	Day of month; valid range 1-31				
Month	21	4	Month of year; 1= January , valid range 1-12				
Year	25	7	Count of years from 1980; Valid range is: 0 — year 1980 127 — year 2107				

The time stamp format used in the ExFAT is same as FAT



## **Stream Extension Directory Entry**

	Stream Extension Directory Entry of ExFAT					
Name	Offset (byte)	Size (bytes)	Description			
Entry Type	0	1	0xC0			
General Secondary Flags	1	1	Bit 0 : Allocation possible 0 - No cluster allocated; 1 - cluster allocation is possible  Bit 1 : No FAT chain 0 - Yes; The clusters of this file/directory are NOT contiguous 1 - No; The Contiguous Cluster are allocated to this file/directory; This improves the File read performance  Bits 2 - 7 : Reserved			
Reserved	2	1				
Name length	3	1	Length of file name in bytes ( Maximum of 255 Unicode characters)			
Name hash	4	2	Hash of the file name; Used while searching for fie/directory name			
Reserved	6	2				
Valid data length	8	8	The size of the file of directory in bytes			
Reserved	16	4				
First cluster	20	4	First data cluster of the file/directory			
Data length	24	8	The size of the file of directory in bytes; In case of directory maximum size is 256MB			



#### **File Name Extension Directory Entry**

	File Name Extension Directory Entry of ExFAT					
Name	Offset (byte)	Size (bytes)	Description			
Entry Type	0	1	0xC1			
General Secondary Flags	1	1	Bit 0 : Allocation possible 0 - No cluster allocated; 1 - cluster allocation is possible  Bit 1 : No FAT chain 0 - Yes ; The clusters of this file/directory are NOT contiguous 1 - No; The Contiguous Cluster are allocated to this file/directory; This improves the File read performance  Bits 2 - 7 : Reserved			
File name	2	30	15 Unicode character of the part of the file name			

- Single File name Extension directory entry can contain 15 Unicode characters.
- Maximum File name size is 255 Unicode characters
- Hence, for a single long file/directory name maximum of 17 file name extension directory entries are possible.

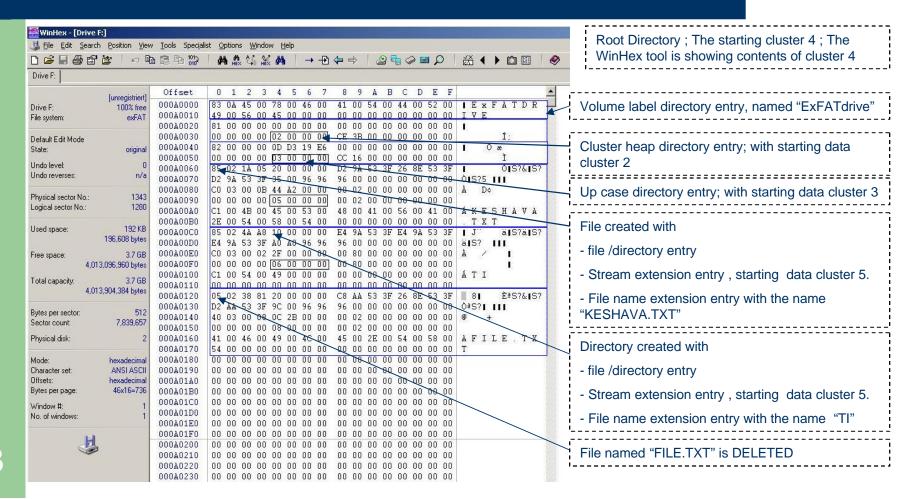


# File/Directory and Volume label Deletion process

- The Most Significant Bit (MSB) of Entry Type field of the directory entry will be changed from 1 to 0 to indicate the directory entry is deleted.
- For example, A file contains the following directory entries
  - File Directory Entry , Entry Type is 0x85 = 0x10000101 in binary format
  - Stream Extension Directory Entry, Entry type is 0xC0 = 0x11000000 in binary format
  - File name extension Directory Entry, Entry type is 0xC1 = 0x11000001 in Binary format
- The MSB (7<sup>th</sup>) Bit indicates "In Use" field; it means 1 In use, 0 deleted.
- This MSB bit is changed from 1 to 0 to indicate the deletion process;
  - File Directory Entry, Entry Type is 0x85 is changed to 0x05 = 0x00000101 in binary format
  - Stream Extension Directory Entry, Entry type is 0xC0 is changed to 0x40 = 0x01000000 in binary format
  - File name extension Directory Entry, Entry type is 0xC1 is changed to 0x41= 0x010000001 in binary format
- The Entry type of the volume label is 0x83 = 0x10000011 in binary format.
- This Entry type value 0x83 will changed to 0x03 = 0x00000011 to indicate that there is no volume label.

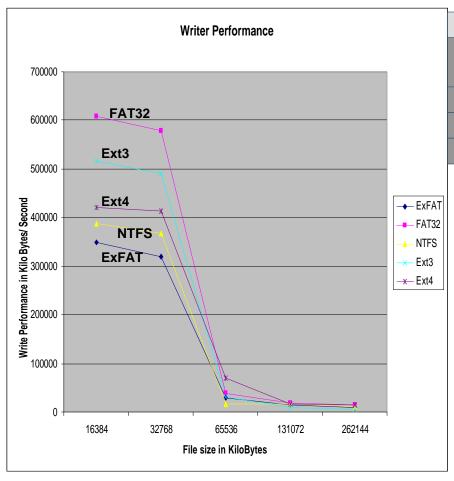


### **Snapshot of the Root Directory**





# File systems Write performance numbers in Linux

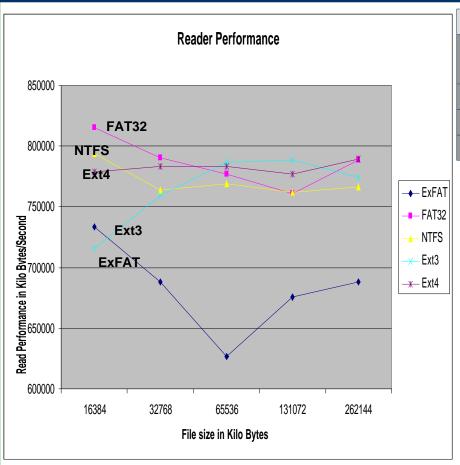


Test Setup		
Computer	Dell Computer, Optiplex GX620	
	Duel core 3.4Ghz with 2GB RAM	
Operating system	Ubuntu 9.04 with Linux kernel version 2.6.28	
Storage device	Transcend 4GB Thumb drive	
Performance tool	lozone version 3.3	

- The fie size range 16MB to 256MB
- The buffer size: 4MB to write
- For small files higher performance range is 300MB/S (Mega bytes per second) to 700MB/S.
- For larger files, the performance range is 1MB/S to 12MB/S.
- The ExFAT is implemented as FUSE( File system in User space) module.
- FAT32 has seen good and in performance for both smaller and larger files.



# File systems Read performance numbers in Linux

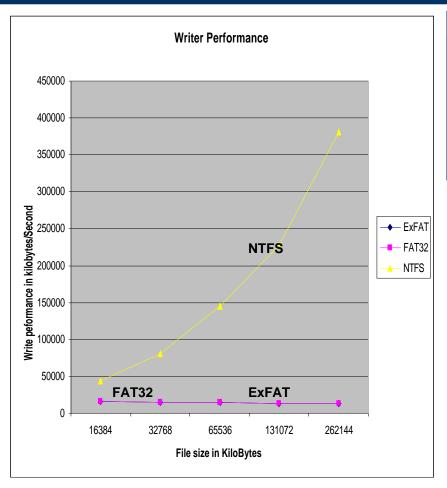


Test Setup	
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	Duel core 3.4Ghz with 2GB RAM
Operating system	Ubuntu 9.04 with Linux kernel version 2.6.28
Storage device	Transcend 4GB Thumb drive
Performance tool	lozone version 3.3

- The fie size range 16MB to 256MB
- The buffer size: 4MB to Read
- For small files higher performance range is 700MB/S (Mega bytes per second) to 850MB/S.
- For larger files, the performance range is 600MB/S to 800MB/S.
- Compared to write, the read performance degradation is less.
- FAT32 and Ext4 has seen good and consistent in performance for both smaller and larger files.



# File systems Write performance numbers in WindowsXP

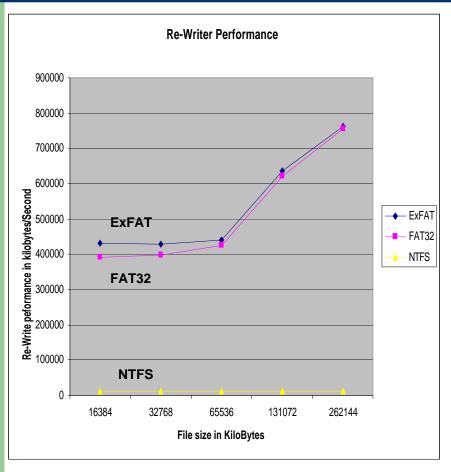


Test Setup		
Computer	Dell Lap-top, Latitude D630	
	Intel Core 2Duo CPU 2.4 GHz, &1.17Ghz	
Operating system	Windows XP Profession version 2002, Service pack3	
Storage device	Transcend 4GB Thumb drive	
Performance tool	lozone version 3.3	

- The fie size range 16MB to 256MB
- The buffer size: 4MB to Write
- NTFS is has highest performance range 50MB/S to 400MB/S
- Both ExFAT and FAT32 has lower performance less than 50 MB/s
- The NTFS performance is higher because of "Write caching" mechanism
- In Window XP, user has to set the "Write caching mechanism by enabling the performance optimization policy in the Hardware tab of the drive.



# File systems Re-Write performance numbers in WindowsXP

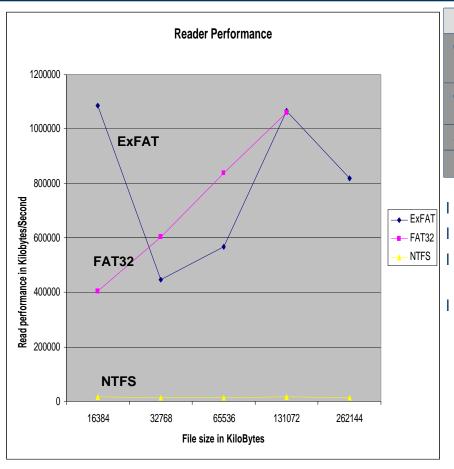


Test Setup		
Computer	Dell Lap-top, Latitude D630	
	Intel Core 2Duo CPU 2.4 GHz, &1.17Ghz	
Operating system	Windows XP Profession version 2002, Service pack3	
Storage device	Transcend 4GB Thumb drive	
Performance tool	lozone version 3.3	

- Re-Write means write the file that already exists.
- The fie size range 16MB to 256MB
- The buffer size: 4MB to Write
- NTFS is has lowest performance below 100 MB/S
- Both ExFAT and FAT32 has higher performance range 400MB/s to 800MB/s
- ExFAT has slightly higher performance than FAT32.



# File systems Read performance numbers in WindowsXP



Test Setup		
Computer	Dell Lap-top, Latitude D630	
	Intel Core 2Duo CPU 2.4 GHz, &1.17Ghz	
Operating system	Windows XP Profession version 2002, Service pack3	
Storage device	Transcend 4GB Thumb drive	
Performance tool	lozone version 3.3	

The fie size range 16MB to 256MB
The buffer size: 4MB to Read
NTFS is has lowest performance
below 100 MB/S

Both ExFAT and FAT32 has higher performance range 400MB/s to 2GB(Giga Bytes)/S



#### References

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- FAT Specification: <a href="http://msdn.microsoft.com/en-us/windows/hardware/gg463080">http://msdn.microsoft.com/en-us/windows/hardware/gg463080</a>



#### Questions

#### **Queries and Feedback**

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