

## MINOR - 1 (COL-100)

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## PROBLEM 3:

To Show That :  $\text{divide}(x, y) = (q, r)$ where  $x = qy + r$ for  $x \geq 0, y > 0, q \geq 0, 0 \leq r < y$ Proof : By PMI Version 3 on  $x$ Basis : For  $x=0$   $\text{divide}(x, y)$  returns  $(q, r) = (0, 0)$ IH : For  $0 \leq k < x$ ,  $\text{divide}(k, y)$  returns  $(q, r)$   
Such that  $k = qy + r$   $q \geq 0$  and  $0 \leq r < y$ 

IS :

Case 1 ( $x = 2k$ ) : EVEN

\* Let  $x$  be even i.e.  $x = 2k$ . Then by  
IH  $\text{divide}(x \text{ div } 2, y)$  returns  $(q, r)$  such that  
 $k = qy + r$   $q \geq 0$  and  $0 \leq r < y$ . Thus by  
algorithm

$$x = 2k = q_1 y + r_1$$

$$\text{where } q_1 = 2q \text{ \& } r_1 = 2r$$

If  $r_1 < y$  the algo. returns  $(q_1, r_1)$  as required.

If  $r_1 \geq y$  then  $r_1 - y = 2r - y < y$  (as  $r < y$ ).

Thus  $x = (q_1 + 1)y + (r_1 - y)$  and the algo  
returns  $(q_1 + 1, r_1 - y)$  as required.



Case 2 : ODD ( $x = 2k+1$ )

By IH

divide ( $x \text{ div } 2, y$ ) returns  $(q, r)$

$$k = qy + r, \quad q \geq 0, \quad 0 \leq r < y$$

Thus by algo

$$x = 2k+1 = q_1 y + r_1 + 1 = q_1 y + r_2$$

$$\text{where } q_1 = 2q, \quad r_2 = 2r+1$$

If  $r_2 \leq y$  algo returns  $(q_1, r_2)$  as required.

If  $r_2 > y$  then

$$r_2 - y = 2r+1 - y < y \quad \text{because } r < y$$

Thus  $x = (q_1+1)y + (r_2 - y)$  and the algo

returns  $(q_1+1, r_2 - y)$  as required.

Hence Proved by PMI  $\forall 3$

# Finding Time Complexity

$$x = 2^m$$

$$T(x) = T(x \text{ div } 2) + O(1)$$

$$= m \cdot O(1)$$

$$= T(2^m \text{ div } 2) + O(1)$$

$$= m \cdot O(1)$$

$$= k \log_2 x$$

$$\therefore T(x) = O(\log_2 x)$$