8086 Memory Organizations

Physical Address / Real Address

SEGMENT: OFFSET CNCEPT

The number of Address Bits depends on the amount of memory we want to handle, at maximum, a 16-bit processor can access memory of up to $64KB~(2^{16})$. If we want to handle memories of more than 64KB we need additional bits.

In the 8086, all registers are <u>16 bits</u> wide, but the address bus requires 20-bit wide. In order for an address of <u>20 bits</u> in width to be composed in code, the CPU automatically combines the contents of two registers for all memory access. The <u>two</u> registers that combine to make up the 20-bit address are called generally <u>Segment Registers</u> and Offset Registers.

There are four Segment Registers (CS, SS, DS, and ES), and numerous other Offset Registers (AX, BX, CX, DX, IP, SP, BP, SI, & DI), and they combine in ways, to compose the 20-bit address.

Some Basic Concepts:

1. Physical Address:

It is the 20-bit address that actually put on the address bus (in 8086). It has a range of 00000_h – FFFFF_h. The physical address or real address refers to the actual position in the memory. It allows access to data in the main memory.

2. Logical Address:

A logical address or virtual address is used as a reference to access the physical address. Because the physical address cannot be accessed directly. The logical address consists of segment address and offset address.

Logical address = Segment addr : Offset addr

3. Segment Address:

It is a 16-bit address of the segment block. The segment address, located within one of the segment registers, defines the beginning address of any 64K-byte memory segment.

4. • Offset Address :

It is a location within 64K byte segment range. It has a range of 0000h - FFFFh

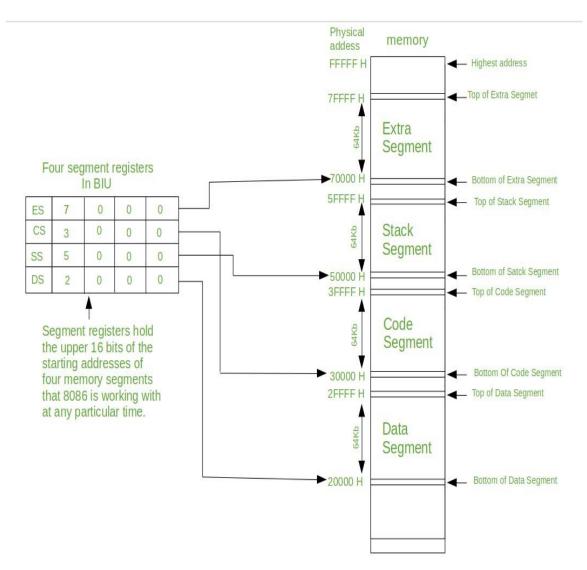


Fig. Shows 8086 Memory Segmentation

To express the 20-bit Physical-address of memory, Follow the steps below:

1. Multiply Segment register by 10h (or shift it to left by 4 bit).

2. Add it to the offset. (See Figure below)

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This is performed automatically by the EU or CPU, and does not require any code to perform the address arithmetic.

Segment address indicates Segment Number, while \underline{Offset} address indicates the location of byte or word \underline{within} the segment.

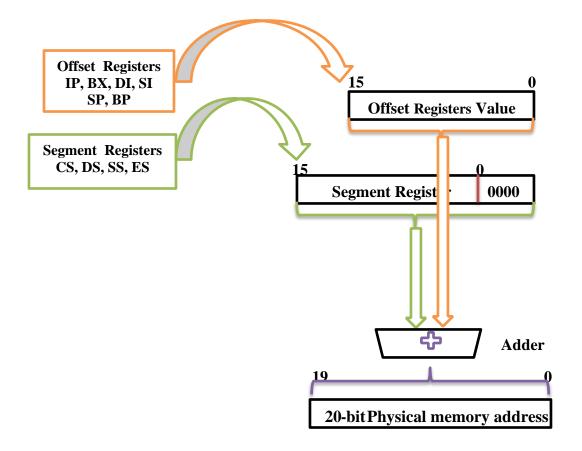


Figure Shows Physical Address And How To Calculate

Addressing in Code Segment:

To execute a program, the 8086 fetches the instructions from the Code Segment. The logical address of an instruction *consists* <u>CS</u> (Code Segment) and <u>IP</u>(instruction pointer). Logical Address in Code segment is *represented* by using segment address in CS register and Offset Address in IP register as follows:

General Formula of logical address:

Segment addr.: Offset addr.

Segment Address 16 bit CS: IP Offset Address 16 bit

Example:

Assume, CS register contains 2500h and IP register contains 95F3h.

What is the Logical Address in the Code Segment?

CS:IP
$$\rightarrow$$
 2500:95F3

The 2500:95F3 means offset of 95F3 within Segment 2500

Similarly, the \underline{SS} (Stack Segment) register combines with the \underline{SP} (Stack Pointer) Offset Register to access the stack.

SS:SP or SS:BP

Physical Address / Actual Address

The following operation illustrates the generation of 20-bit physical address from segment: offset pair 2222: 118C

<u>Sol</u>:

Multiply Segment register by 10h (or shift it to left by 4 bit). 22220

Add it to the <u>offset</u>. 118C

physical address 20-bit 233AC

Example:

Calculate from segment: offset pair, If SS = 24F6h and SP = 634Ah, Show the:

1- Logical Address SS:SP = 24F6: 634A

2- Offset Address SP = 634A

3- Physical Address = 2B2AA physical address 20-bit

4- The Lower range of the Segment 24F6:0000 => 24F60 + 0000 = 24F60

5- The Upper range of the Segment 24F6: FFFF = 24F60 + FFFF = 34F5F

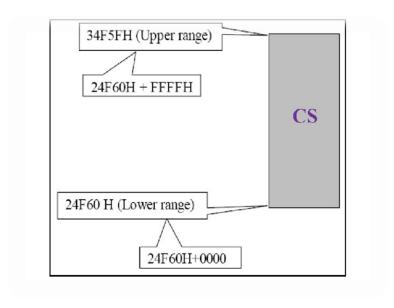


Figure Shows Lower / Upper range of the Segment

Example:

Physical Address = 1278B, Code Segment = 1278, Find the

Offset value & Logical Address?

<u>Sol</u>:

From: 12780 + Offset = 1278B

Offset = 000B

Logical Address = CS: IP = 1278: 000B

Location of Segments:

Segment 0 starts at address 0000:0000 = 00000h and ends at 0000:FFFF = 0FFFFh.

Segment 1 starts at address $0001:0000 = 00010_h$ and ends at $0001:FFFF = 1000F_h$.

Segment 2

Segment 3 etc.