CS241 #15

Condition Variables. Implement a Mutex Lock. The Critical Section Problem

- 1. How do I block a thread (= send it to 'sleep')?
- 2. How do I wake up threads that are blocked on a condition variable?

Example: Fix the following methods using a condition variable and mutex lock to ensure the cake integer is never negative.

```
pthread mutex t m = PTHREAD=MUTEX INITIALIZER;
    pthread cond t cv = PTHREAD COND INITIALIZER;
04
   int cake = 0;
    void decrement() { // Will block if zero
      while(cake == 0) {
10
          sleep(1)
11
12
      cake --;
14
15
16
17
    void increment() {
18
        cake ++;
19
```

3. How does pthread cond wait really work?

4. Challenge. A fixed size stack:

```
01 pthread mutex t m = PTHREAD=MUTEX INITIALIZER;
02 pthread cond t cv = PTHREAD COND INITIALIZER;
03 double array[10];
04 \text{ int } n = 0;
06 // blocks while full (n ==10)
07 void push(double v) {
09
10
11
12
13
14
15
16 }
17 //  blocks while empty (n == 0)
18 double pop() {
19
20
21
22
23
24
25
26 }
27 // Test with 2+ threads that add values...
28 void* generator(void*) {
     for (int i = 0; i < 100000; i++)
         push(i);
    Return NULL;
31
32 }
33 // And one thread that remove values
34 void * consumer(void*result) {
    double sum = 0, i=0;
    while (i=pop() != -1) sum += i;
    printf("%.0f", sum);
    Return NULL;
39 }
40
```

How can you implement a reliable mutex lock?

5. Let's try writing a simple implementation...

```
pthread_mutex_init(int * m) { *m= 0; }

pthread_mutex_lock(int* m) {
    while(*m ==1) {
        pthread_yield(); /*sleeps for a short time */
    }

    *m = 1;

    }

pthread_mutex_unlock(int* m) {?______}
```

Problems?

6. CPU support: Use an atomic CPU instruction.

Suppose a special 'Atomic_Exchange' instruction 'exch' exists that swaps the values at two addresses as an uninterruptable operation

```
01  pthread_mutex_init(int* m) { *m= 0; }
02
03  pthread_mutex_lock(int* m) {
     for(int q = 1; q;) { ______ }
     }
04
05  pthread_mutex_unlock(int* m) { ______ }
```

7. The Critical Section Problem

while(running) {

- 1. Wait to enter the critical section if another thread is in the CS.
- 2. Critical Section Code here. Only one thread in here at a time!
- 3. Leave critical section. Allow another waiting thread to enter.
- 4. // do other stuff most of the time

~~ Welcome to the **Critical Section Problem** game show! ~~ Today's prizes: mutual exclusion and progress

Candidate #1. Use a single, boolean "flag"

boolean flag

Thread A	Thread B
wait while the flag is up	wait while the flag is up
raise the flag!	raise the flag!
Critical Section code here	Critical Section code here
lower the flag!	lower the flag!

// Then each thread does other work but will repeat this again sometime in the future Problems?

Candidate #2. Give each thread its own a flag.

boolean flagA, flagB

wait while B's flag is up	wait while A's flag is up
raise A flag	raise B flag
Critical Section code here	Critical Section code here
lower A flag	lower B flag

Problems?

Candidate #3. Change the sequence order

raise A flag	raise B flag
wait until B flag is down	wait until A flag is down
Critical Section code here	Critical Section code here
lower A flag	lower B flag

Problems?

Candidate #4. Try a single turn-based shared variable.

turn=1

	while(turn == 1) { }
Critical Section code here	Critical Section code here
turn = 2	turn = 1

Problems?