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Today: Build Your Memory Allocator

typedef struct \_metadata\_entry\_t {

void \*ptr;

int size;

int free;

struct \_metadata\_entry\_t \*next;

} metadata\_entry\_t;

metadata\_entry\_t \*metadata = NULL;

// TA simple solution. i) Complete the code. ii) Which placement algorithm does this use? iii) Does this implementation use explicit or implicit linked list? How would you change this to use a first-fit placement allocation? iv) Why does this code suffer from false fragmentation?

void \*malloc(size\_t size) {

/\* See if we have free space of enough size. \*/

metadata\_entry\_t \*p = metadata;

metadata\_entry\_t \*chosen = NULL;

while (p != NULL) {

if (p->free && \_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_) {

if (chosen == NULL || (chosen && p->size < chosen->size)) {

chosen = p;

}

}

p = p->next;

}

void free(void \*ptr) {

// "If a null pointer is passed //as argument, no action occurs."

if (!ptr)

return;

// Free the memory in our metadata.

metadata\_entry\_t \*p = metadata;

while (p) {

if (p->ptr == ptr) {

p->free = 1;

return;

}

p = p->next;

}

return;

}

if (chosen) {

chosen->free = 0;

return chosen->ptr;

}

/\* Add our entry to the metadata \*/

chosen = sbrk(0);

sbrk(sizeof(metadata\_entry\_t));

chosen->ptr = sbrk(0);

if (sbrk(size) == (void\*)-1) {

return NULL;

}

chosen->size = size;

chosen->free = 0;

chosen->next = metadata;

metadata = chosen;

return chosen->ptr;

}

Implementation – Key Ideas

Placement algorithm. Given a linked list of free spaces which one shall I use?

Natural Alignment : Platform able to store all standard C primitives at that address. Platform specific but it is typical: malloc(..) % 16 ==0

External Fragmentation: When the available space is not contiguous. Depends on pattern of allocations and frees.

vs

Internal Fragmentation: ‘Hidden unused space’ inside each allocation   
(standard example: round up each allocation request to 2^n => unused space *inside* each block)

**Implementation notes:**

1. Implicit linked list: Store size of block and calculate offsets to next block

-> Solving Coalescing Problem “False Fragmentation? Use Knuth73 "Boundary Tags" so we can coalesce backwards too.

O(N) alloc. O(1) free.

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| 100 | .................... | 100 |
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| 100 | .................... | 100 | 64 | ....... | 64 | 132 | ............................... | 132 |
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| 100 | .................... | 100 | 64 | ....... | 64 | 132 | ............................... | 132 |
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free(void\*userp) {

if(!userp) return;

myentry\_t\* start = userp - sizeof(myentry\_t);

myentry\_t\* start = ((myentry\_t\*)userp) - 1;

myentry\_t\* prevEnd = start -1;

...

}

2. Explicit linked list: Store memory addresses of next link

Store free blocks pointers inside the unused space of the free block.

Free Block list can now be in arbitrary order.

3. Segregated free list: Different lists for different sizes

Advanced implementation ideas:

Buddy Allocator

Slab allocator.

Deferred coalescing