

# Lecture:

# Software-Defined Networks

Presenter: Ben Ujcich

CS 461 / ECE 422

Fall 2019

# About Me

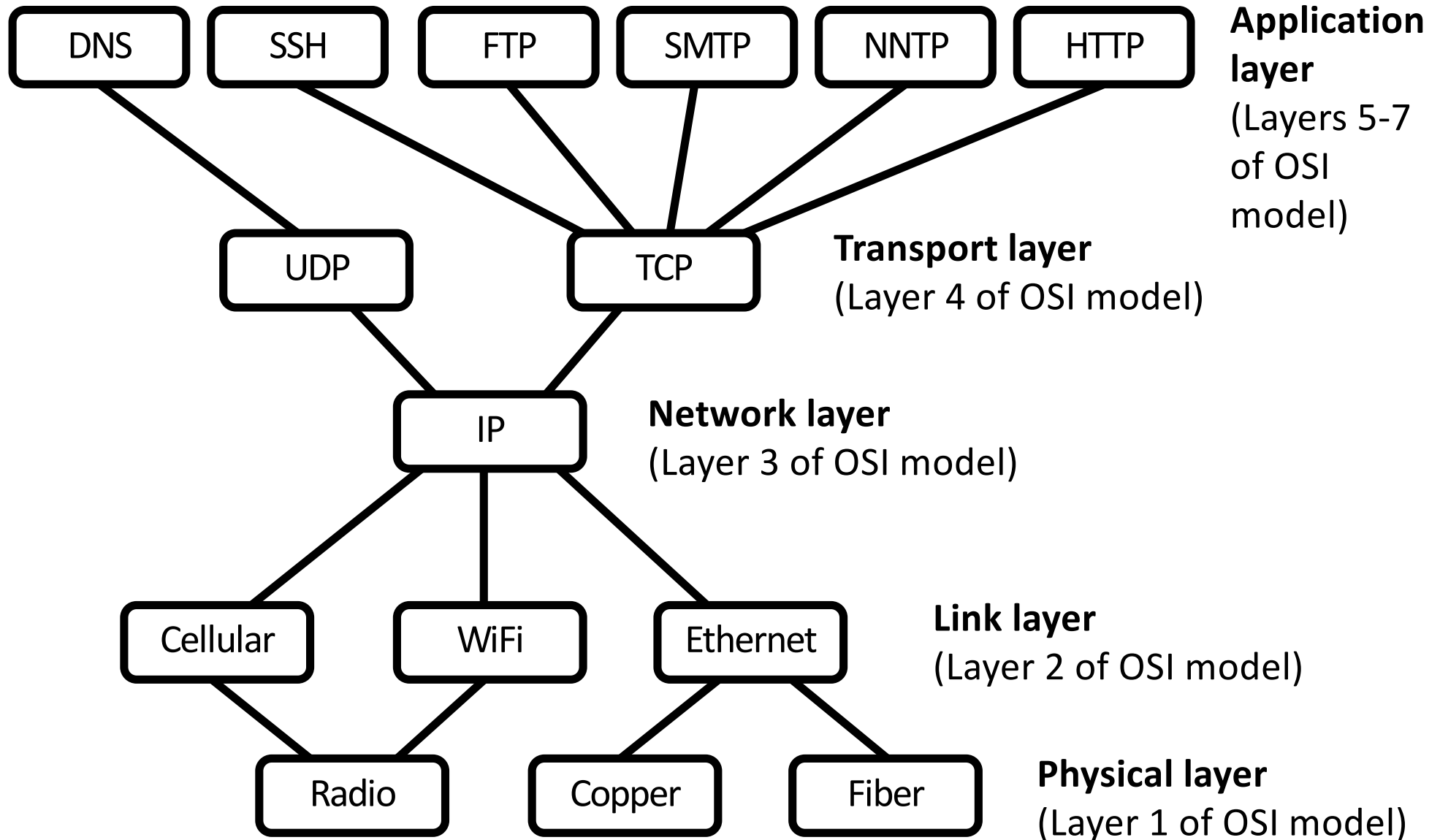
- PhD student in ECE department working with Adam Bates and Bill Sanders
- Broad area of research in security of systems and networks
- Specific area of research in security of software-defined networking (SDN) architectures



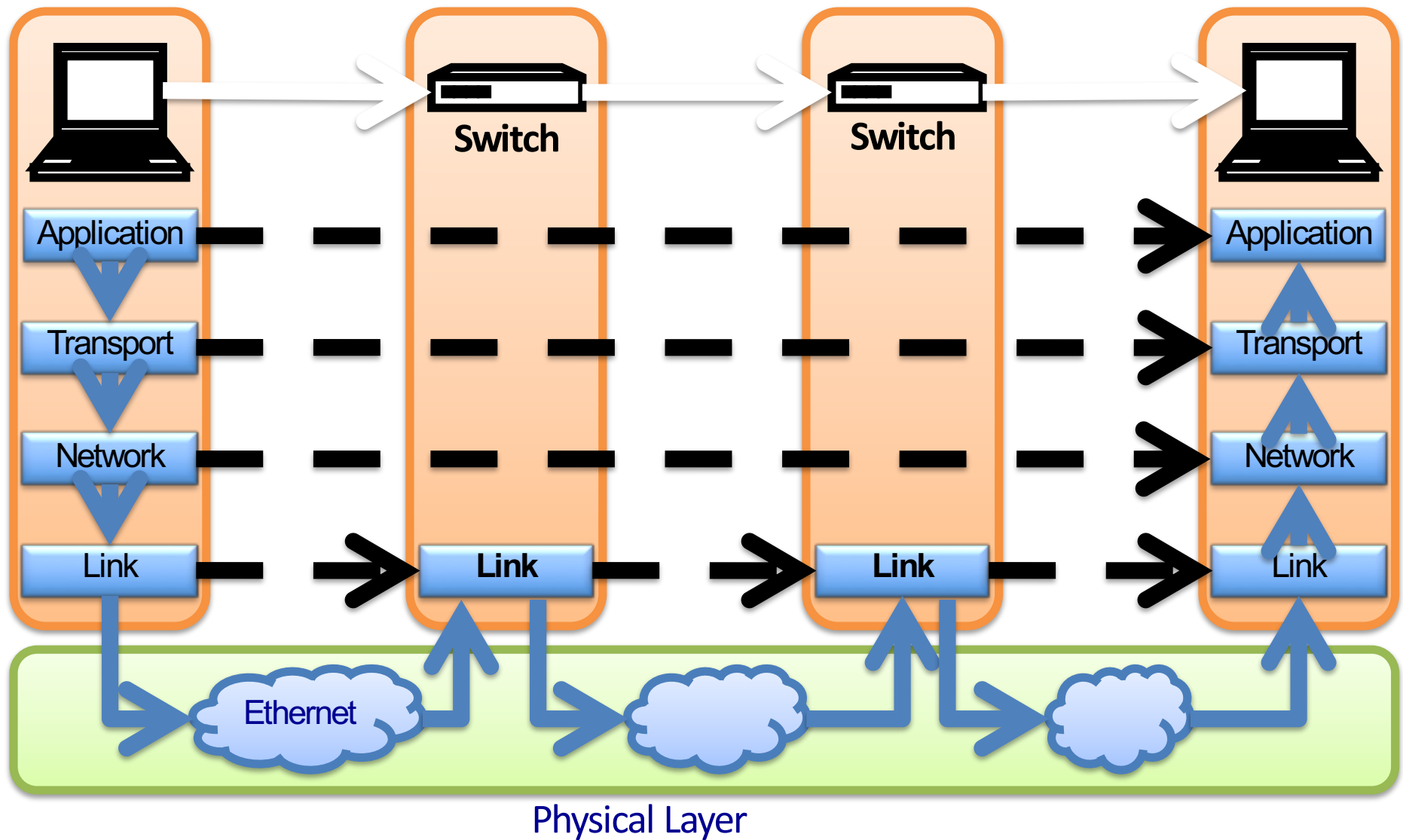
# Goals for Today

- Understand the **challenges** and **limitations** of **traditional networking** architectures
- Understand the **software-defined networking (SDN)** architecture
- Understand the **security benefits** and **ramifications** of SDN
- Understand SDN **attacks** and SDN **defenses**

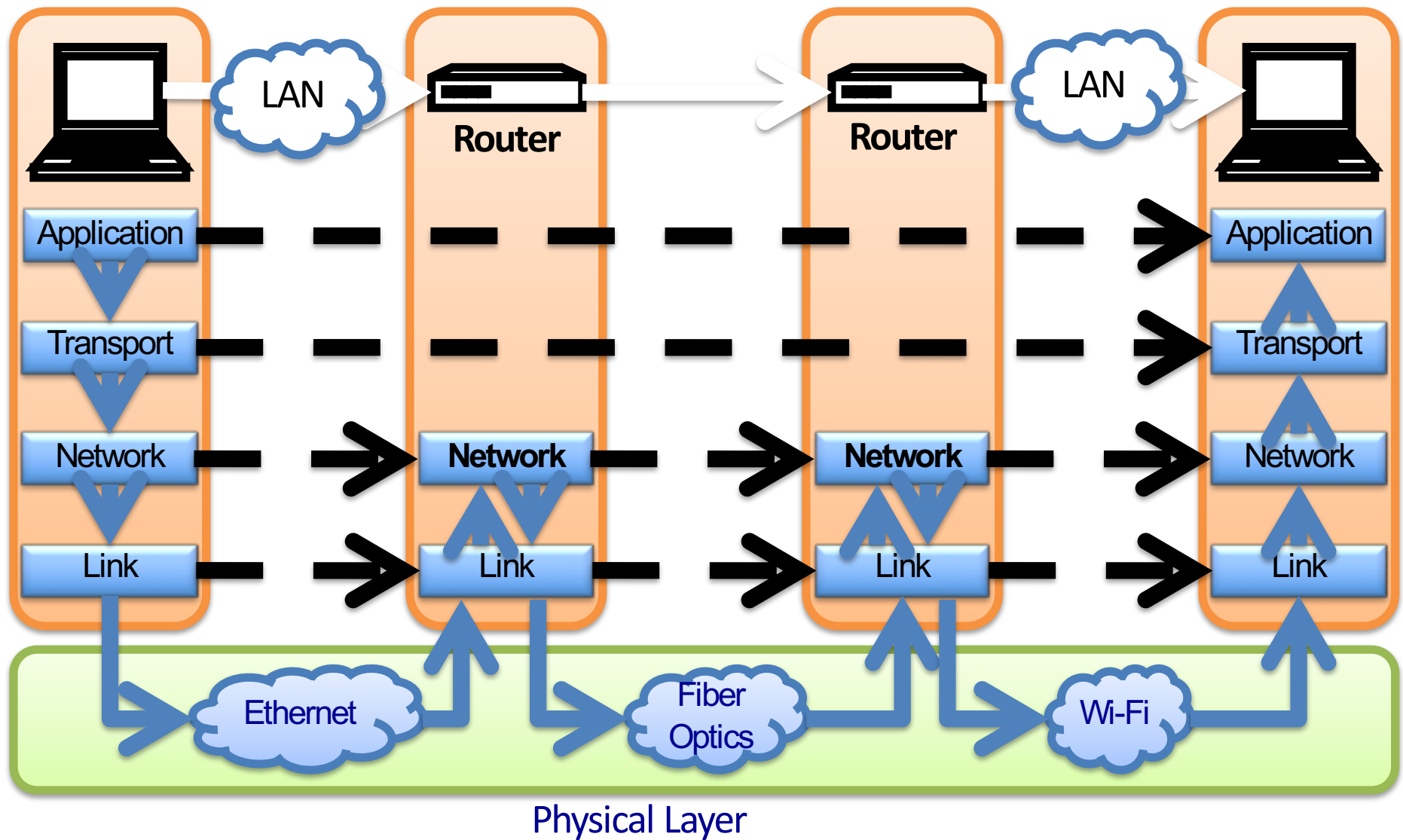
# Recall: Network Stack



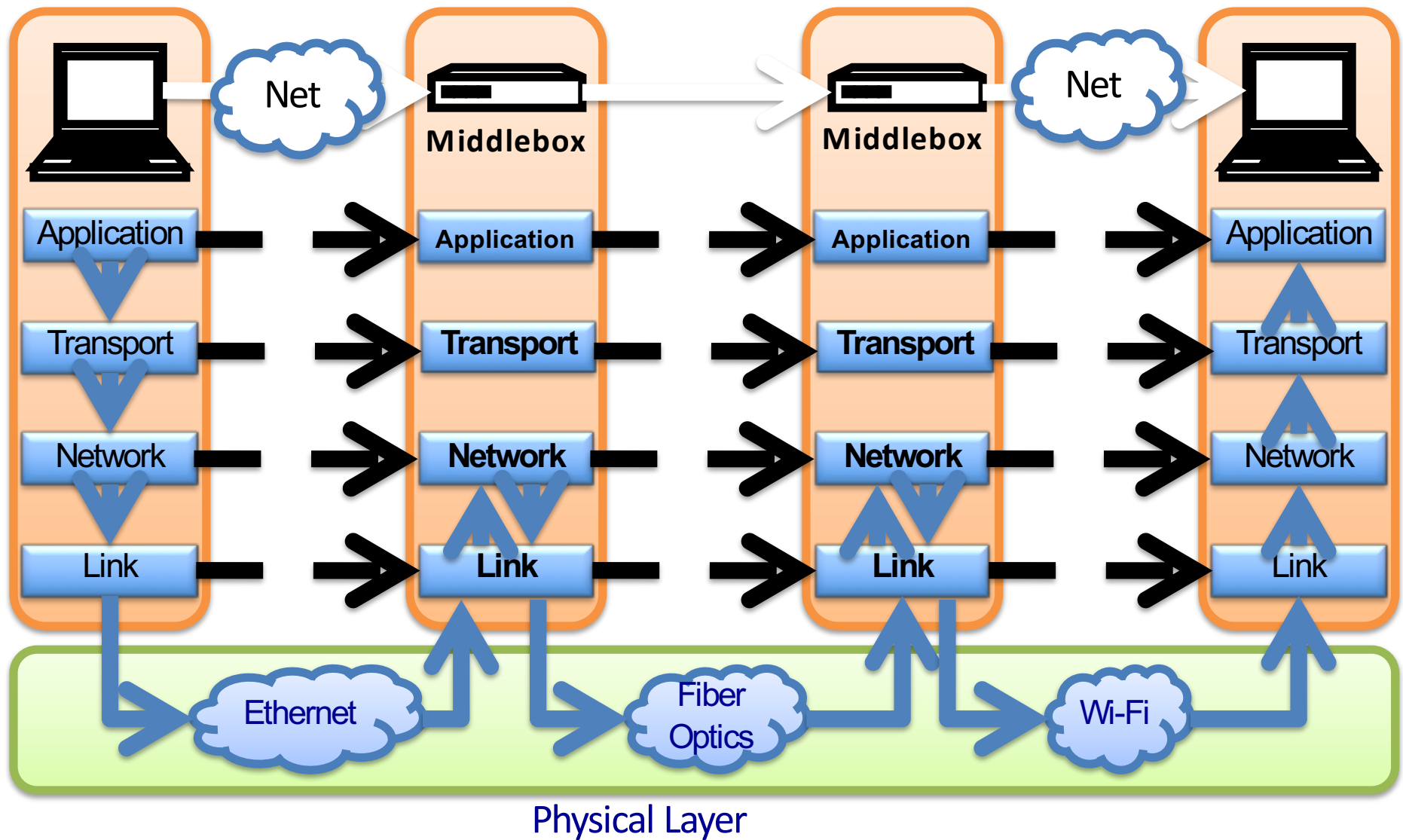
# Recall: Local Area Networks (LANs)



# Recall: Internet



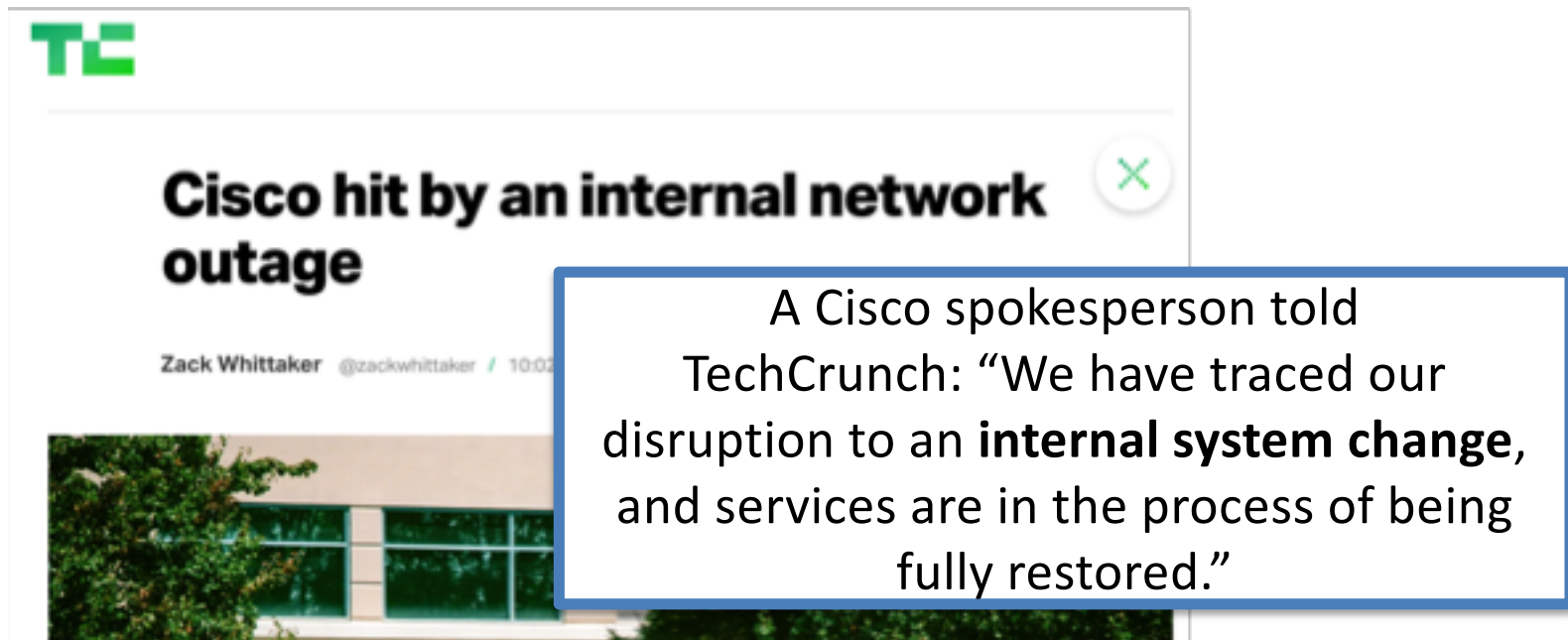
# Higher Layers?



# Challenges

## Difficult to modify network configuration

- Getting the network configuration correct is challenging and error-prone
- Network outages can violate availability property

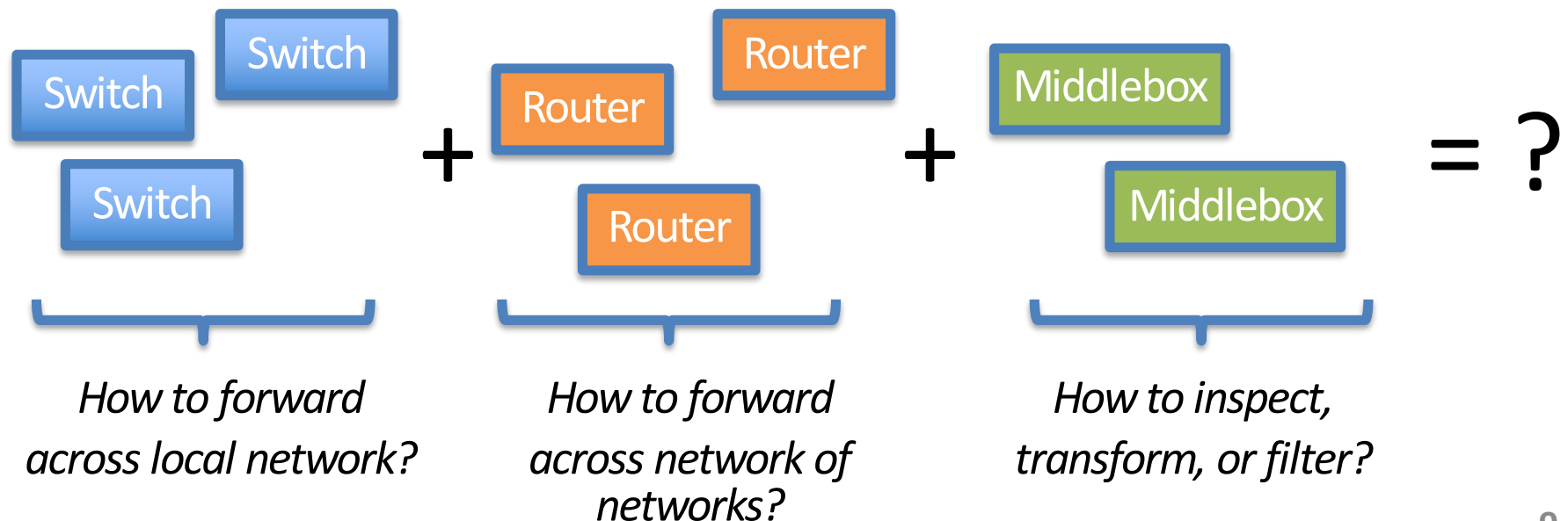




# Challenges

## Difficult to extend / coordinate network functionality and concerns

- Creates cross-layer, cross-protocol problems
- May be limited by physical configuration



# Challenges

## Difficult to separate control and data plane

- The **data plane** forwards packets (or frames or datagrams)
- The **control plane** decides behavior of how packets should be forwarded
- Control and data plane often coupled → *switches care about L2, routers about L3, etc.*
- Individual switches and routers maintain their own views of network state → *network state is distributed*

# Solution: Programmable Networks

- **Insights:**
  1. Decouple control and data planes → *more flexible forwarding across layers of network stack*
  2. Centralize control plane → *coordinated network state and greater insight into decision-making*
  3. Turn decision-making into a set of software processes → *an extensible control plane for an application domain of interest*
- Programmable networks realized as **software-defined networking** or **SDN**

# SDN Data Plane Design

## Traditional Networks

- **Switches** implement forwarding tables based on **MAC addresses**

L2 <sub>src</sub>	L2 <sub>dst</sub>	Action

- **Routers** implement forwarding tables based on **IP addresses / prefixes**

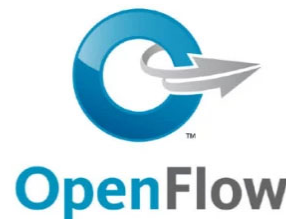
L3 <sub>src</sub>	L3 <sub>dst</sub>	Action

## SDN

- Generic forwarding devices (“switches”) implement forwarding tables based on **Layers 2–4 header fields**

L2 <sub>s</sub>	L2 <sub>d</sub>	L3 <sub>s</sub>	L3 <sub>d</sub>	...	Action

Commonly realized with OpenFlow protocol



# SDN Data Plane Design

## Traditional Networks

- Switches implement forwarding tables based on MAC addresses

**Takeaway:**

**Network stack layers in SDN can be mixed**

- Routers implement forwarding tables based on IP addresses / ranges

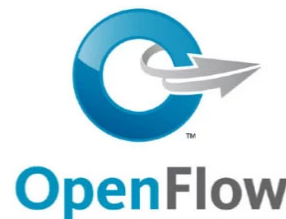
L3 <sub>src</sub>	L3 <sub>dst</sub>	Action

## SDN

- Generic forwarding devices implement forwarding tables based on combination of Layers 2–4 attributes

L2 <sub>s</sub>	L2 <sub>d</sub>	L3 <sub>s</sub>	L3 <sub>d</sub>	...	Action

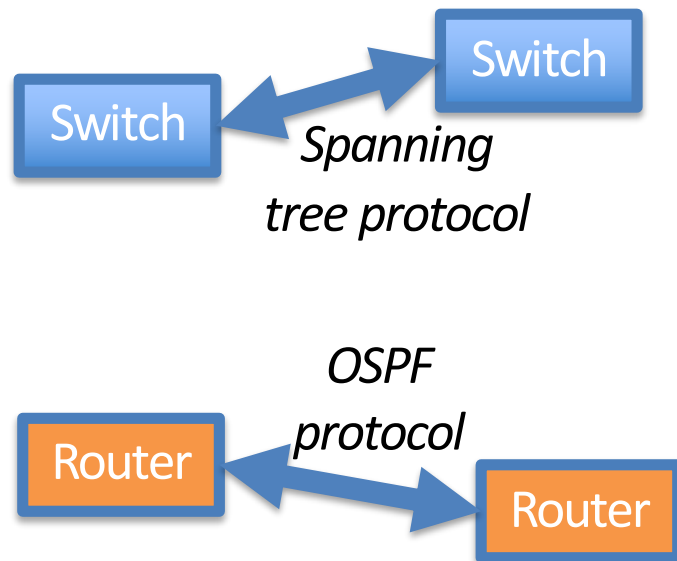
Commonly realized with OpenFlow protocol



# SDN Control Plane Design

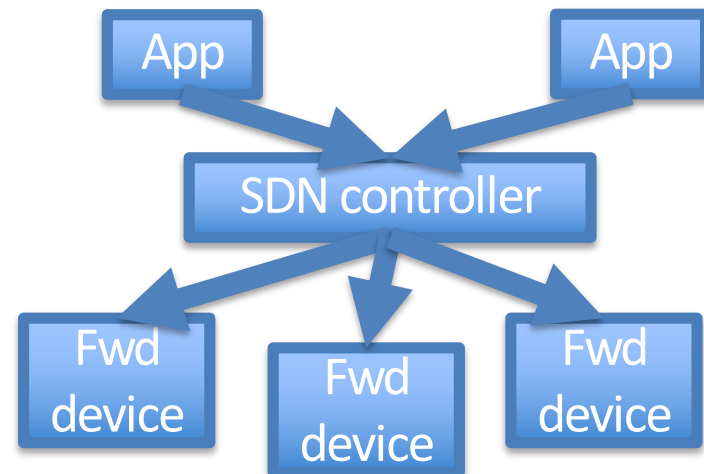
## Traditional Networks

- **Distributed protocols** (e.g., STP, OSPF, BGP, etc.)
- Decision-making resides **on each device**



## SDN

- **Centralized protocols** implemented in software programs (or apps)
- Decision-making resides in **SDN controller** (OpenFlow as configuration protocol)



# SDN Control Plane Design

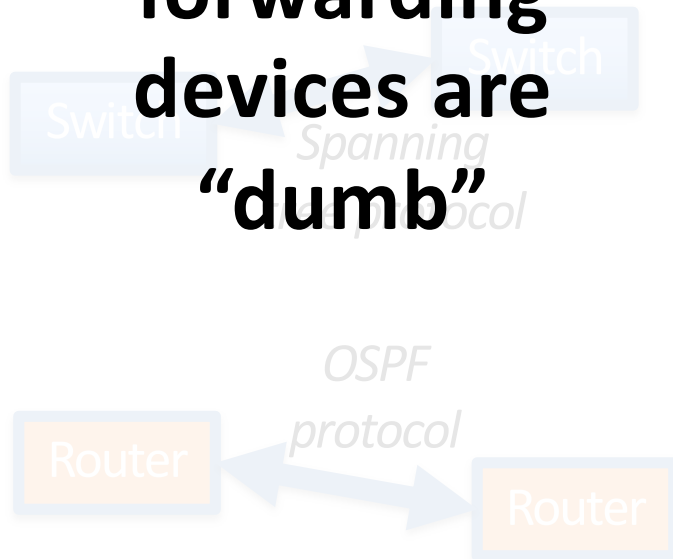
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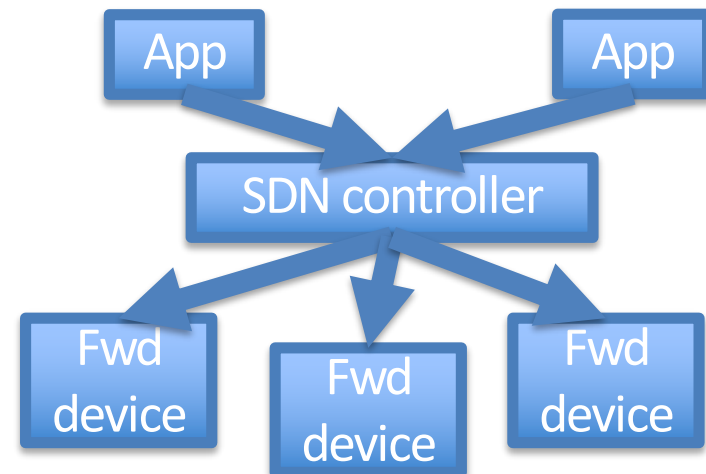
- Decision-making resides on each device

**Takeaway:**  
**Network forwarding devices are “dumb”**

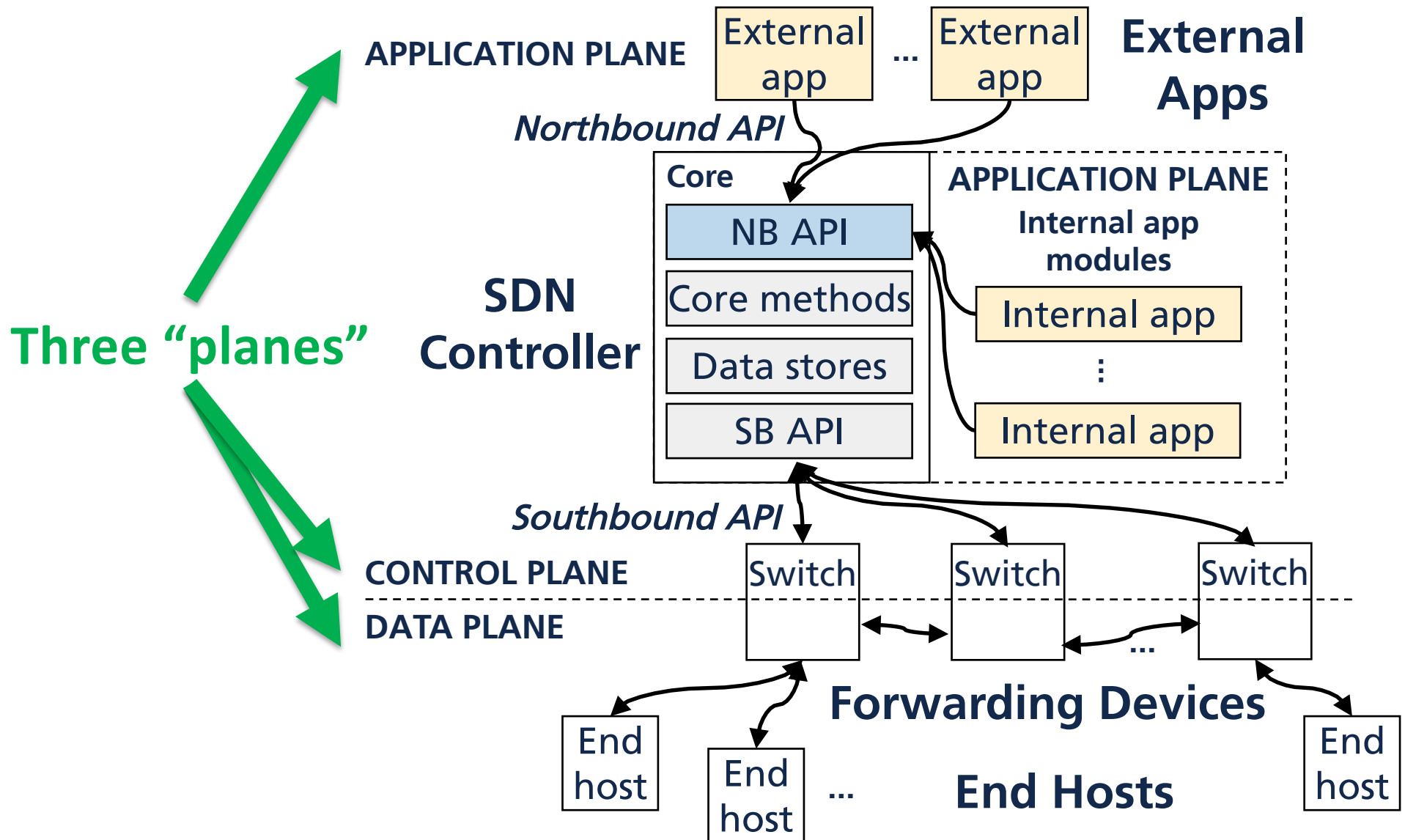


## SDN

- **Centralized protocols** implemented in software programs (or apps)
- Decision-making resides in **SDN controller** (OpenFlow as configuration protocol)

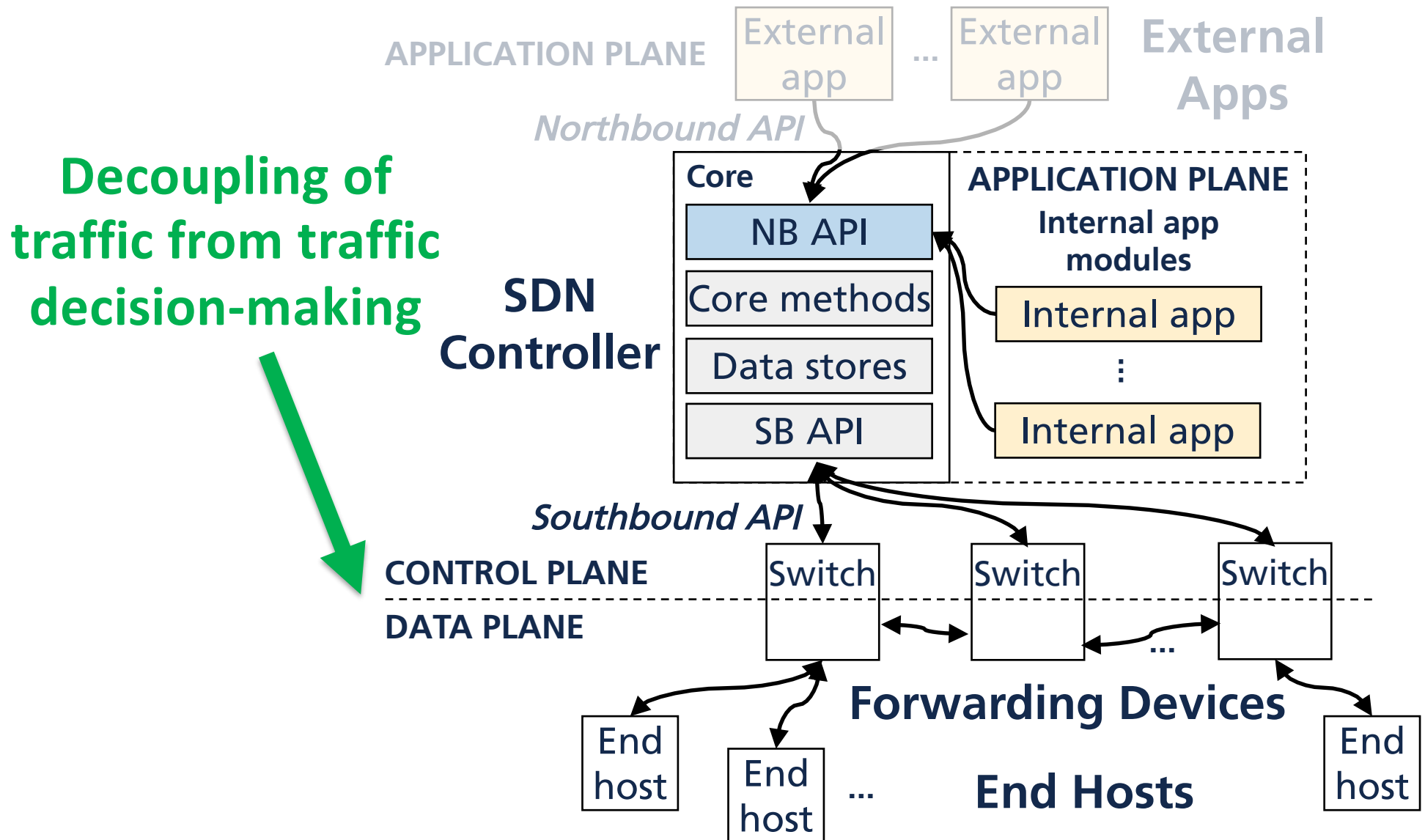


# SDN Architecture

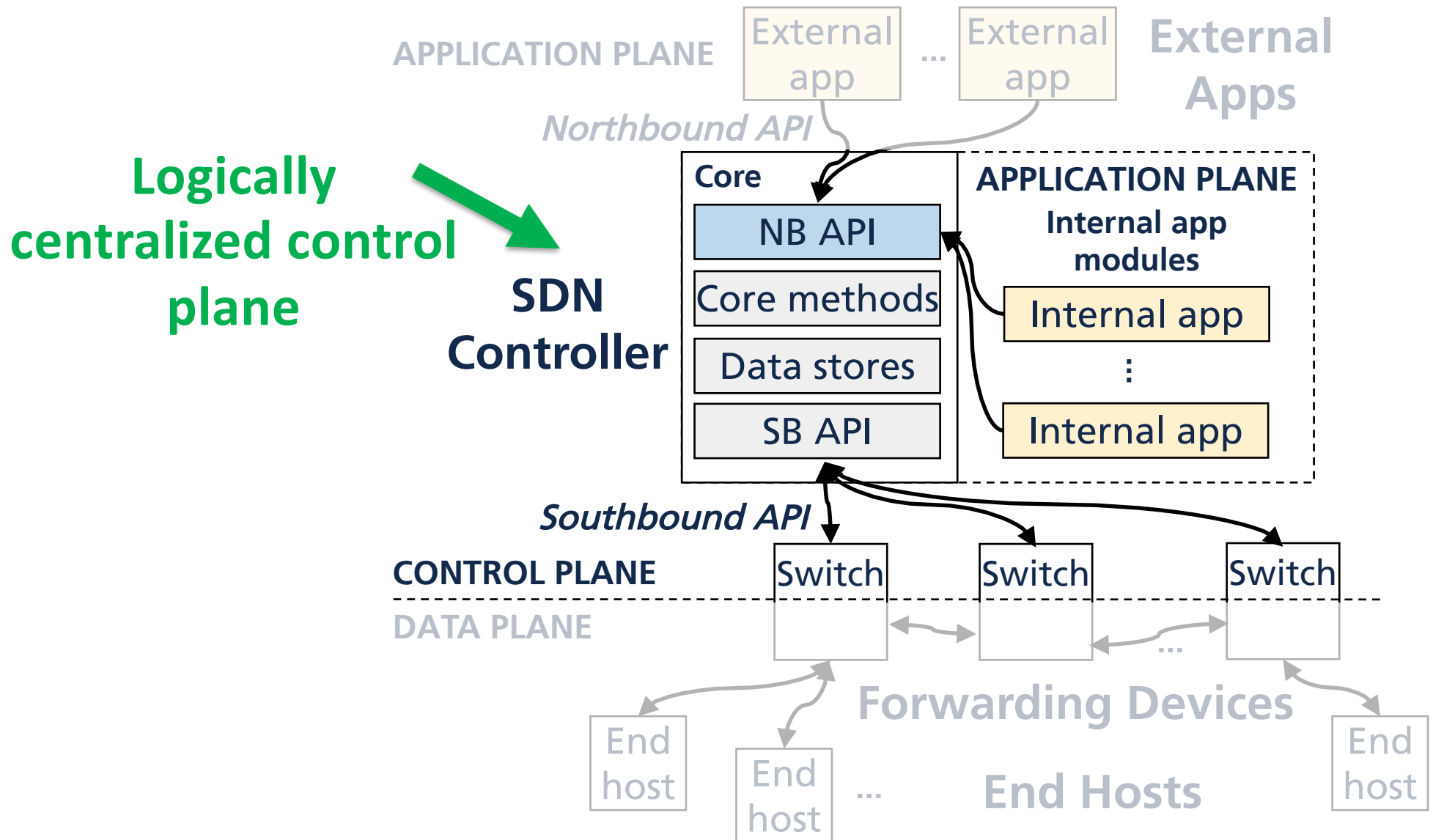




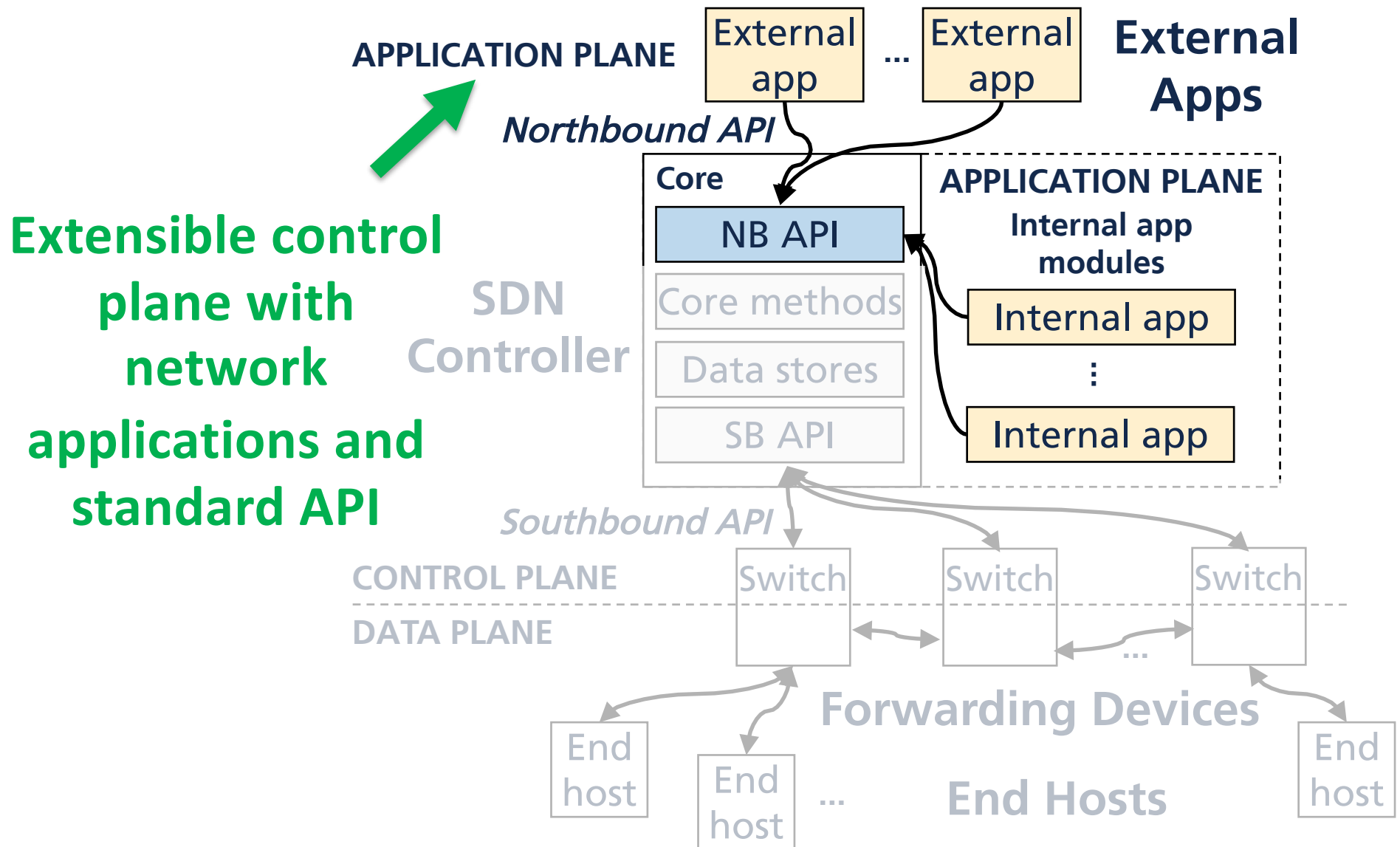
# SDN Architecture



# SDN Architecture



# SDN Architecture



# SDN Controllers and Frameworks

- Open-source
  - Open Network Operating System (ONOS)
  - OpenDaylight (ODL)
  - Floodlight
- Proprietary
  - HPE Virtual Application Networks (VAN) SDN Controller
  - Many other vendor-specific SDNs (e.g., Google B4)



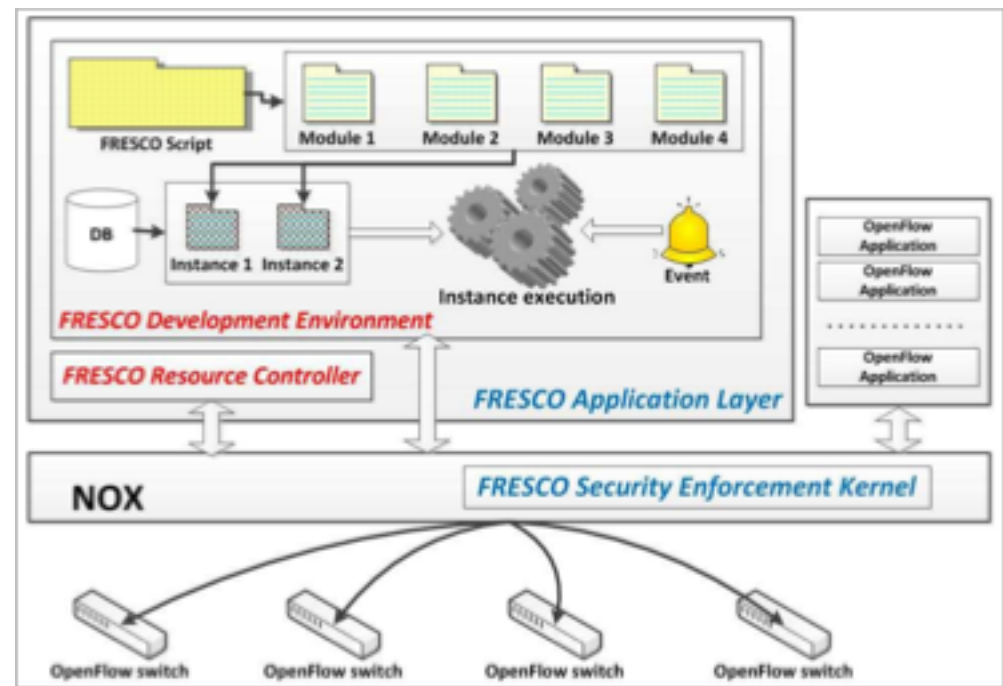
# SDN and Security

- SDN as a security service
  - How can SDN help ensure security properties of the hosts in the data plane?
  - Examples: FRESCO, PSI, DFI
- Security of SDN (attacks and defenses)
  - How does the SDN architecture affect the security properties of the overall network?
  - Examples: Cross-app poisoning, ARP spoofing, LLDP spoofing

# SDN as a Security Service

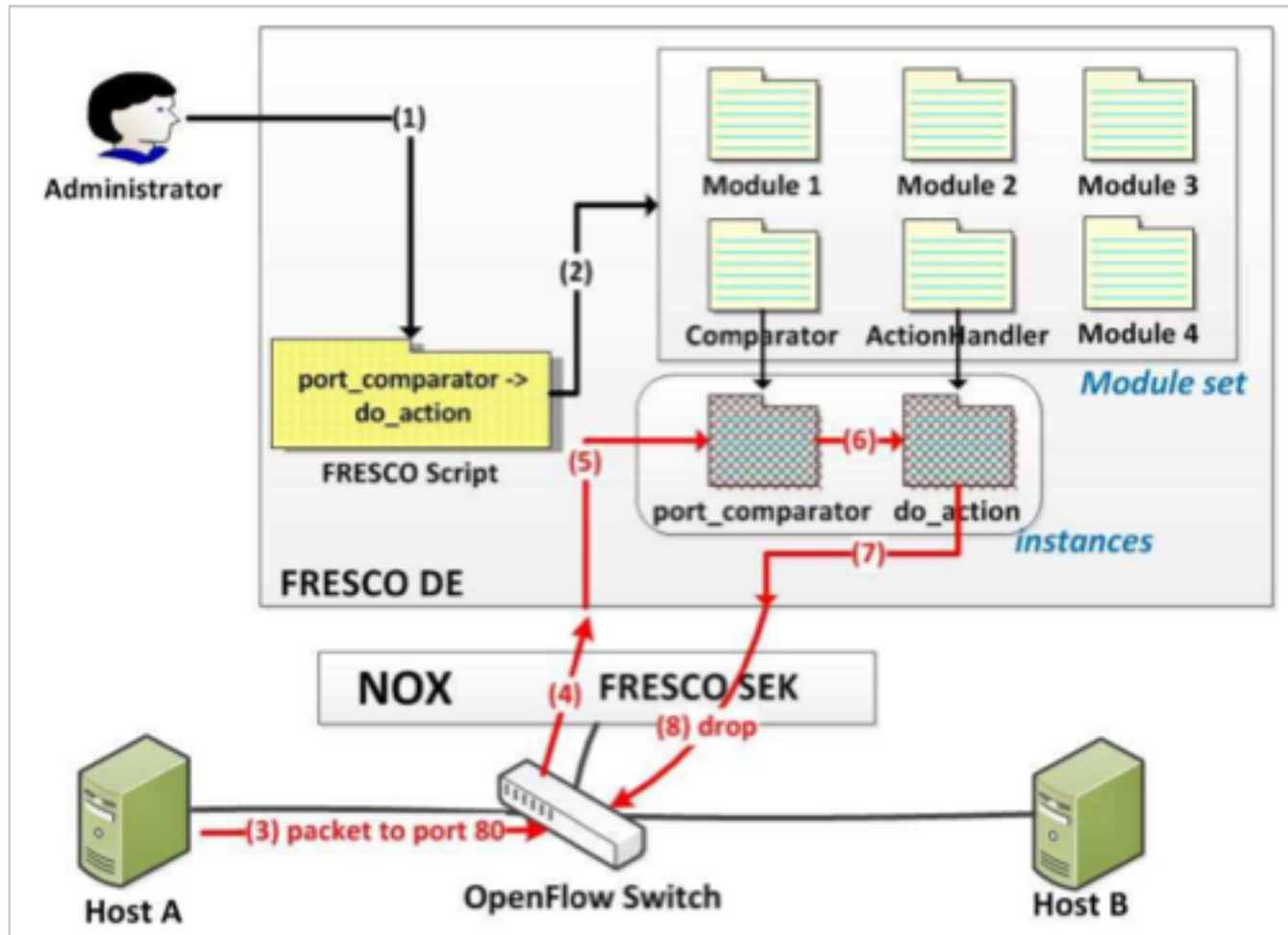
“FRESCO: Modular Composable Security Services for Software-Defined Networks” [NDSS ‘13]

- Problem: Network security applications are difficult to compose correctly
- Solution: Use SDN to develop coordinated security services
- Scripts of modules, inputs, outputs, actions
- Redirect, mirror, quarantine actions



# SDN as a Security Service

“FRESCO: Modular Composable Security Services for Software-Defined Networks” [NDSS ‘13]



# SDN as a Security Service

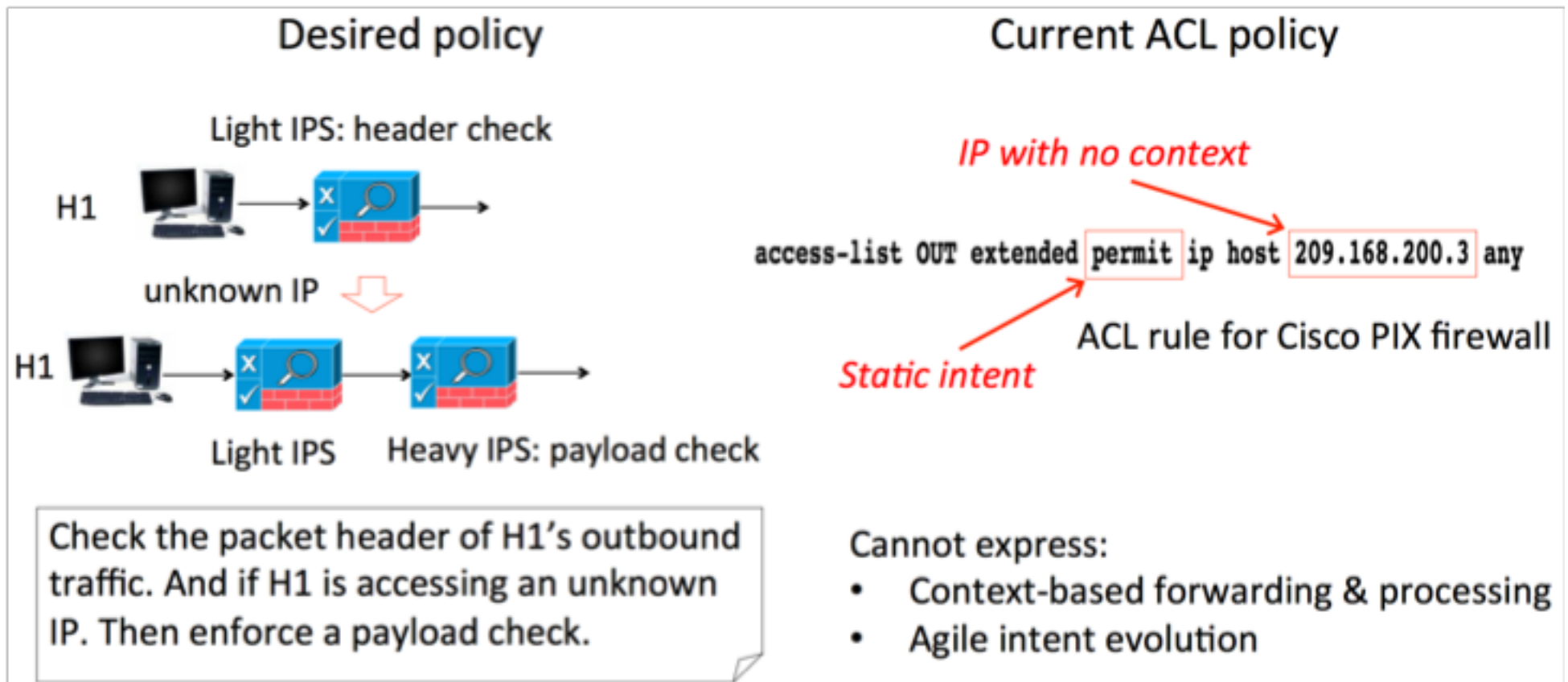
“PSI: Precise Security Instrumentation for Enterprise Networks”  
[NDSS ‘17]

- Problem: Network security suffers from lack of isolation, context awareness, and agility to change rapidly
- Solution: Use SDN to implement network traffic inspection policy with agile reconfiguration
- Enforce policies so that traffic must traverse certain devices (e.g., middleboxes)
- Enforcement policies can be dynamic, so reconfigure network forwarding as needed (e.g., suspiciousness of a potentially compromised host)



# SDN as a Security Service

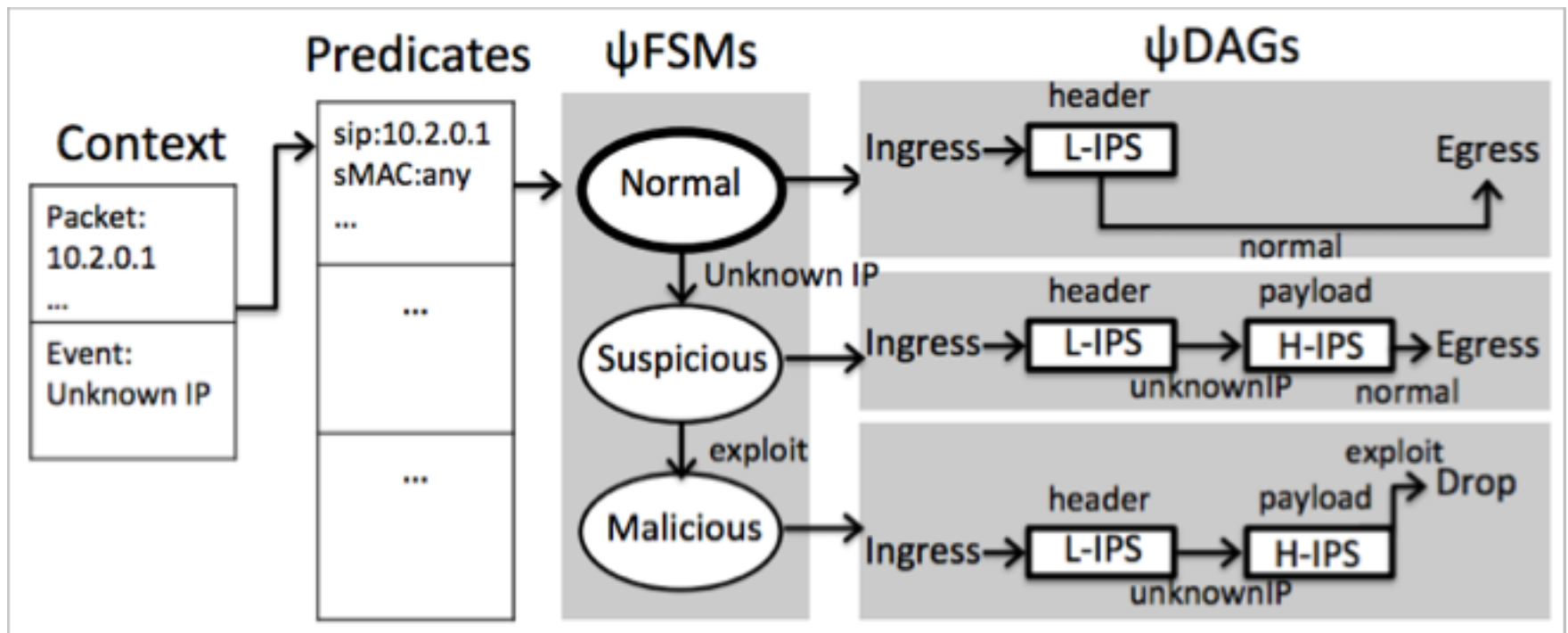
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Source: Slides by authors from paper presentation

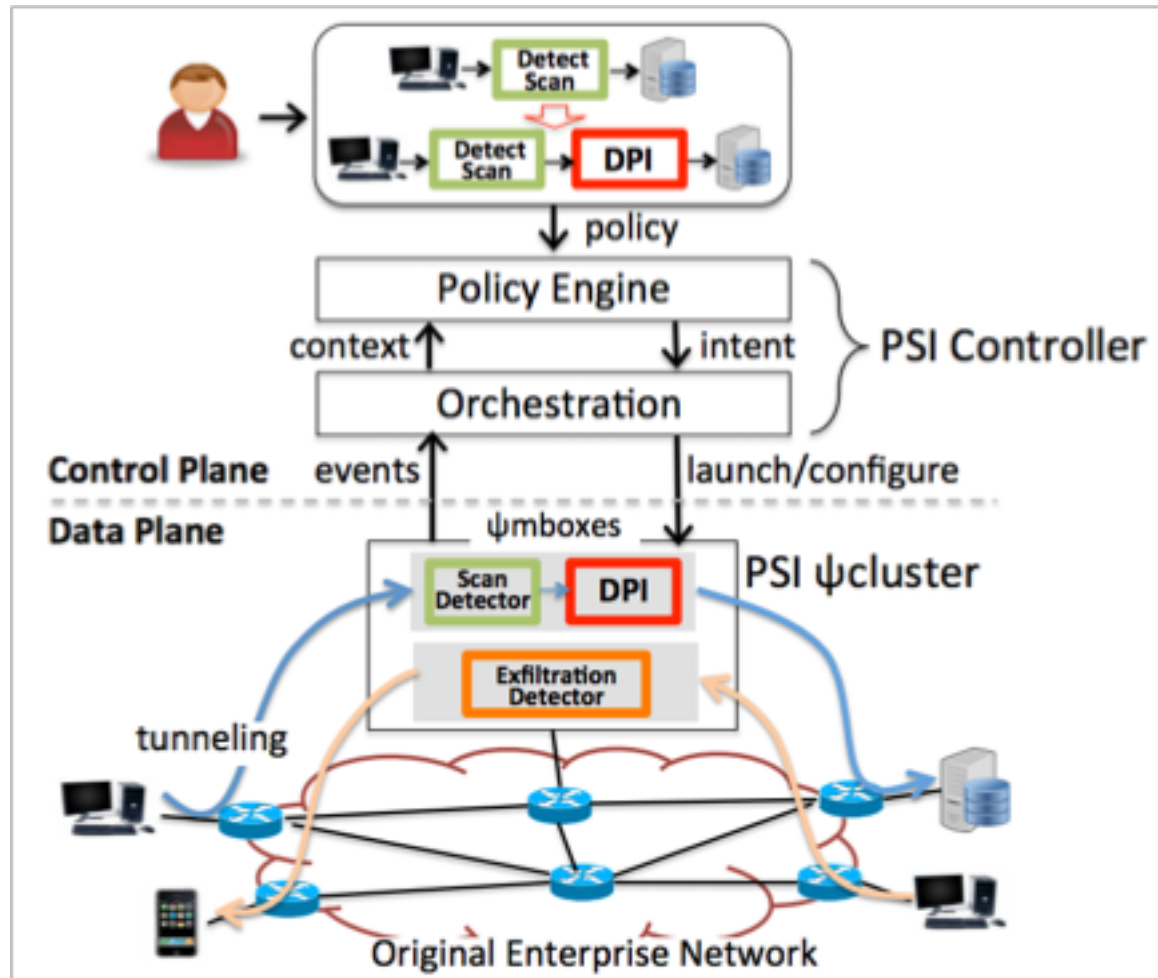
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# SDN as a Security Service

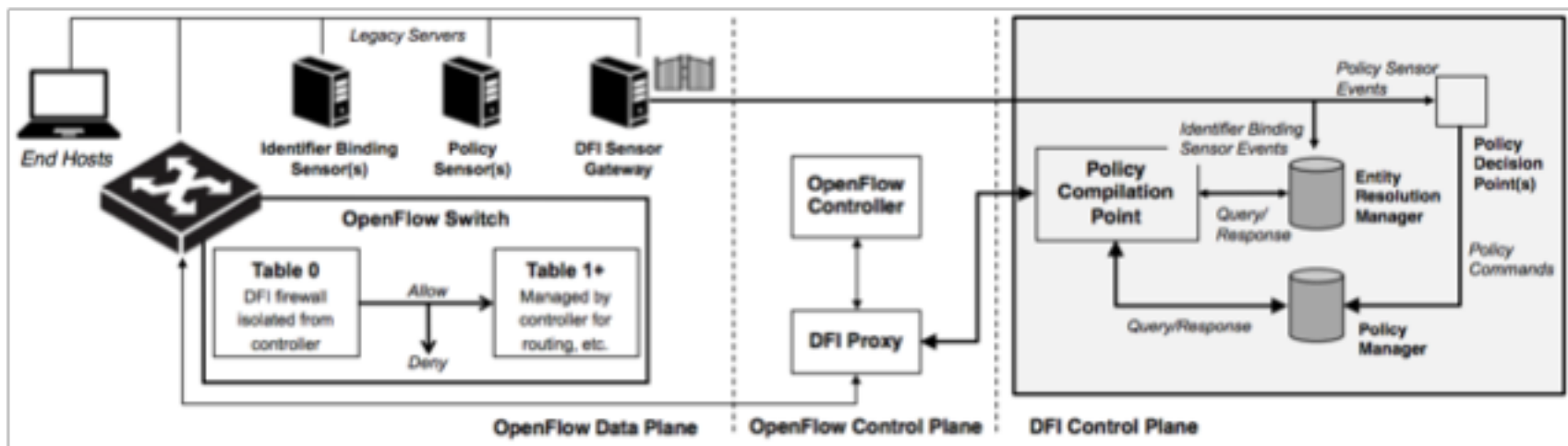
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[NDSS ‘17]



# SDN as a Security Service

“Controller-Oblivious Dynamic Access Control in Software-Defined Networks” [DSN '19]

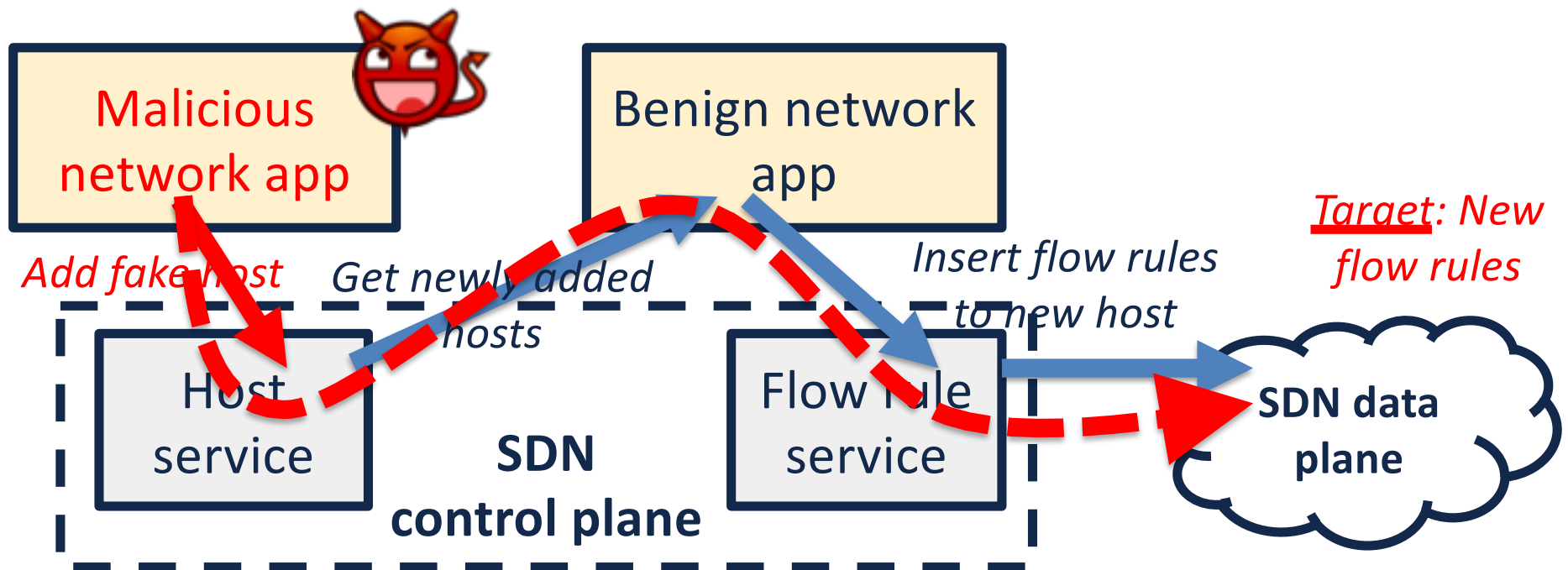
- Problem: Fine-grained, least-privileges access control is necessary to prevent lateral movement, but networks don't respond to non-traditional events (e.g., user logging off)
- Solution: Use SDN and sensors around network to implement access control policies (dynamic flows)



# SDN Attacks: Cross-App Poisoning

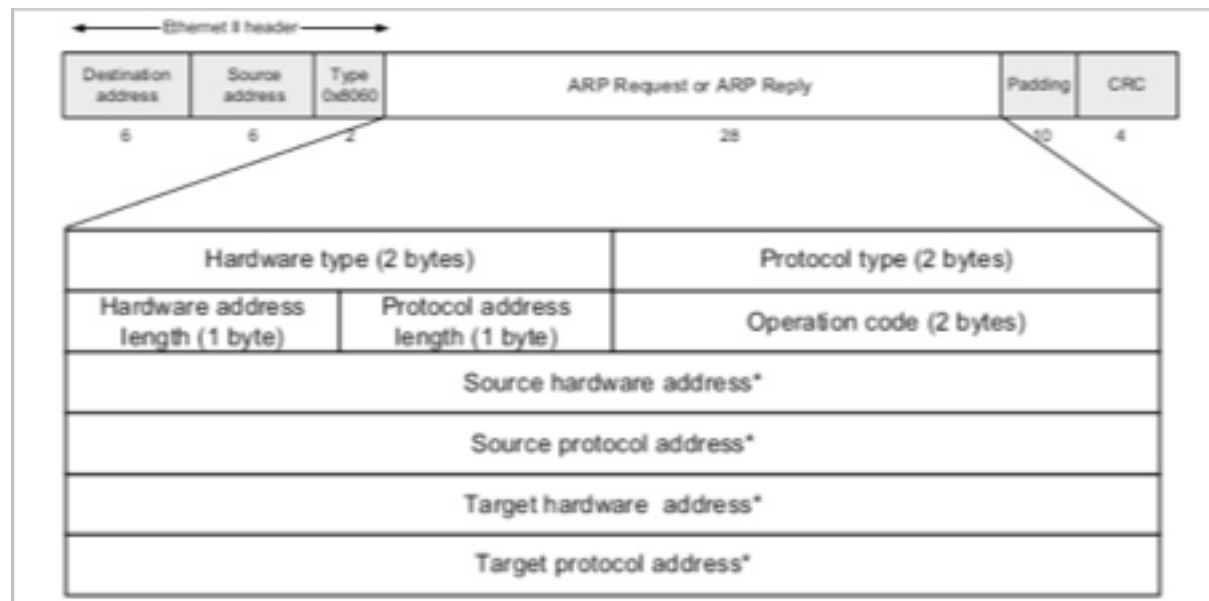
“Cross-App Poisoning in Software-Defined Networking” [CCS ‘18]

- Problem: Malicious network applications can influence other applications through API calls
- Solution: Cross-app poisoning “gadgets” (compare with ROP gadgets); prevent bad information flow



# Recall: Address Resolution Protocol

- Address Resolution Protocol (ARP) lets hosts map IP addresses to MAC addresses
- Host who needs MAC address  $M$  corresponding to IP address  $N$  broadcasts an ARP packet to LAN asking, “who has IP address  $N$ ?”



# Recall: ARP Security

- Any host on the LAN can send ARP requests and replies: *any host can claim to be another host on the local network!*
  - This is called *ARP spoofing*
- This allows any host X to force IP traffic between any two other hosts A and B to flow through X (*MitM!*)
  - Claim  $N_A$  is at attacker's MAC address  $M_X$
  - Claim  $N_B$  is at attacker's MAC address  $M_X$
  - Re-send traffic addressed to  $N_A$  to  $M_A$ , and vice versa

# SDN Attacks: ARP Spoofing

“Poisoning Network Visibility in Software-Defined Networks: New Attacks and Countermeasures” [NDSS ‘15]

- In SDN: the **host location hijacking** attack
- Why this attack works:
  - SDN controllers maintain a database of host objects, derived from data plane packet information
  - As consequence of flattening network stack, controllers use proxy ARP (i.e., single ARP table) → *attacker and victim can be in separate broadcast domains of network*
  - Any data plane host can masquerade as another data plane host
  - No authentication of ARP packets



# SDN Defenses: ARP Spoofing

“Poisoning Network Visibility in Software-Defined Networks: New Attacks and Countermeasures” [NDSS ‘15]

- Static ARP
  - Manually configure ARP table based on known MAC addresses (but hard to change)
- TopoGuard [NDSS ‘15]
  - Uses update semantics (e.g., has host moved?) to mitigate (but not prevent) ARP spoofing
- SecureBinder [USENIX Security ‘17]
  - Use IEEE 802.1x as a root of trust

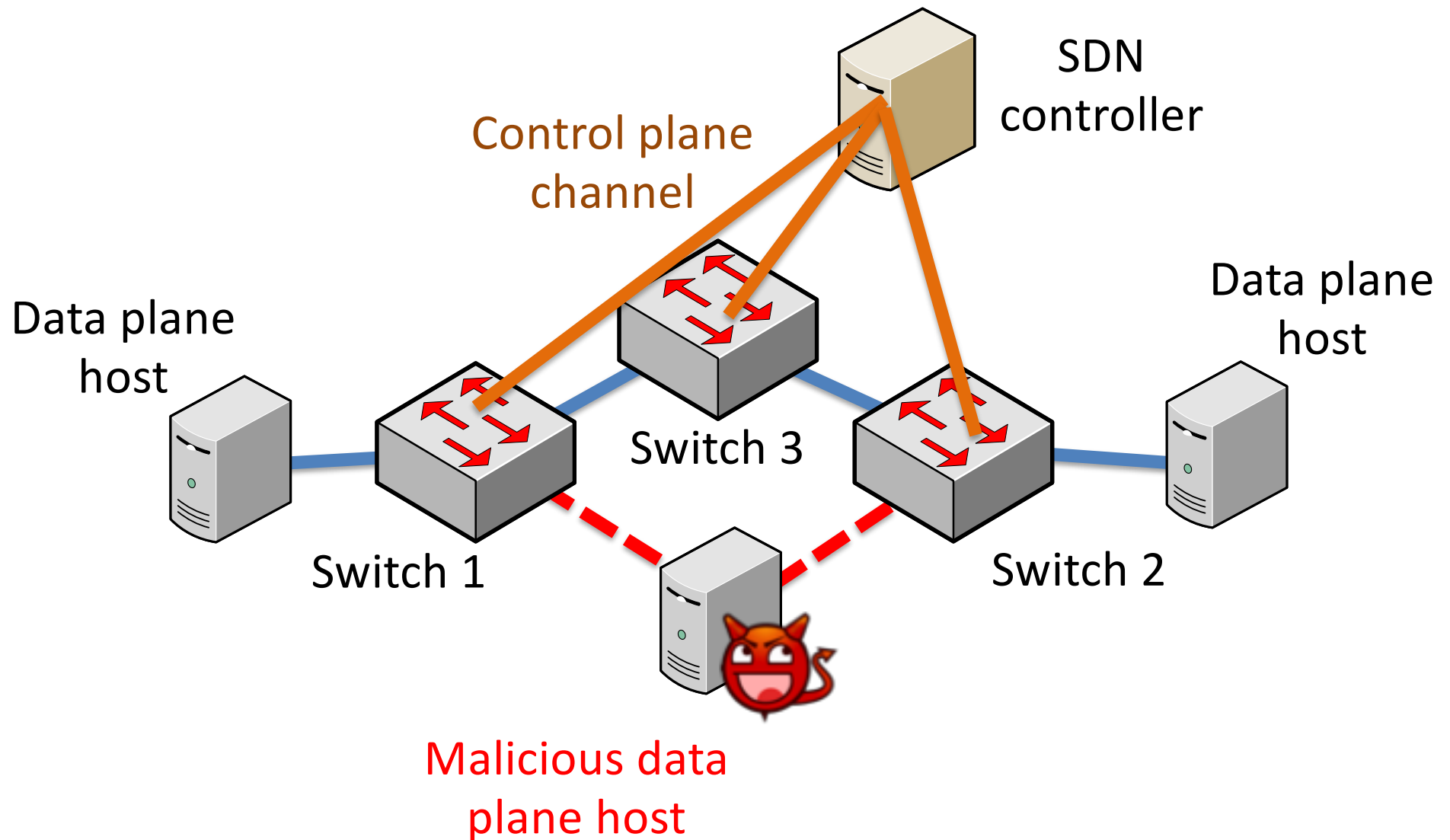
# SDN Attacks: LLDP Spoofing

“Poisoning Network Visibility in Software-Defined Networks: New Attacks and Countermeasures” [NDSS ‘15]

- Link-layer discovery protocol (LLDP) used to learn links (and thus topology) via adjacent switches
- Must use data plane
- Attacker goals:
  1. Create a MitM attack to route traffic through an attacker-controlled host
  2. Create a “black hole” attack to force data plane traffic to be dropped through non-existent links

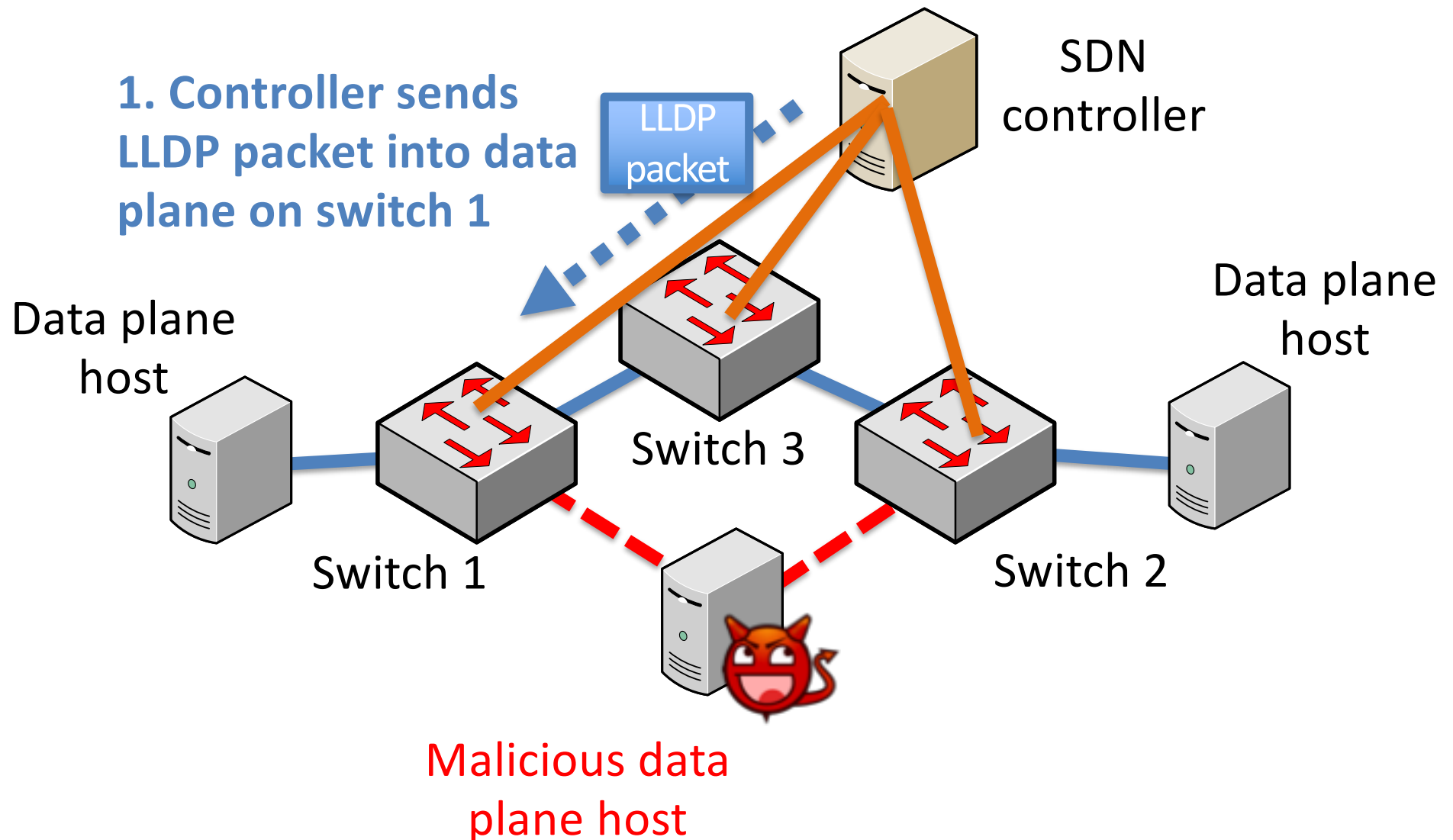
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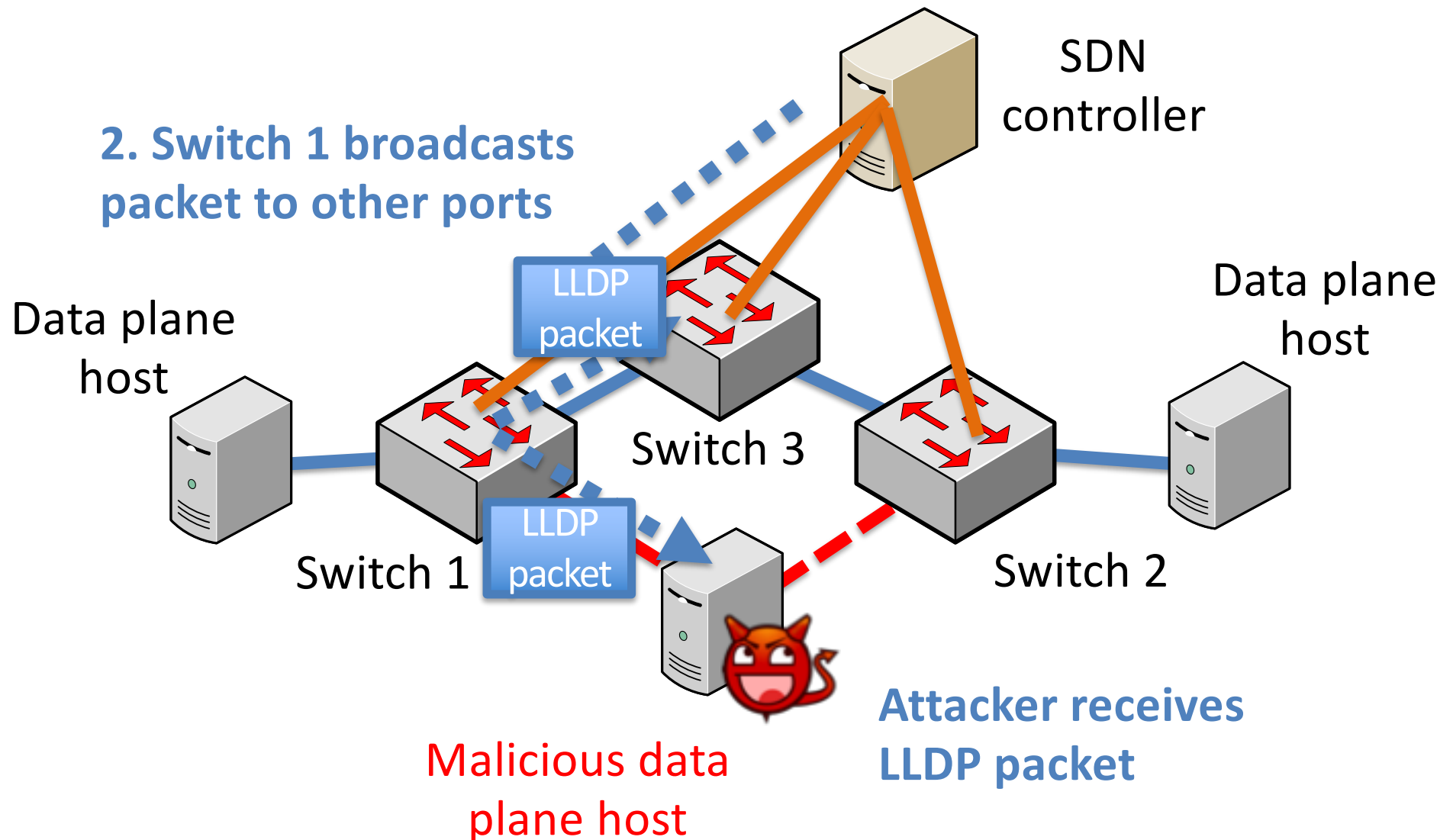
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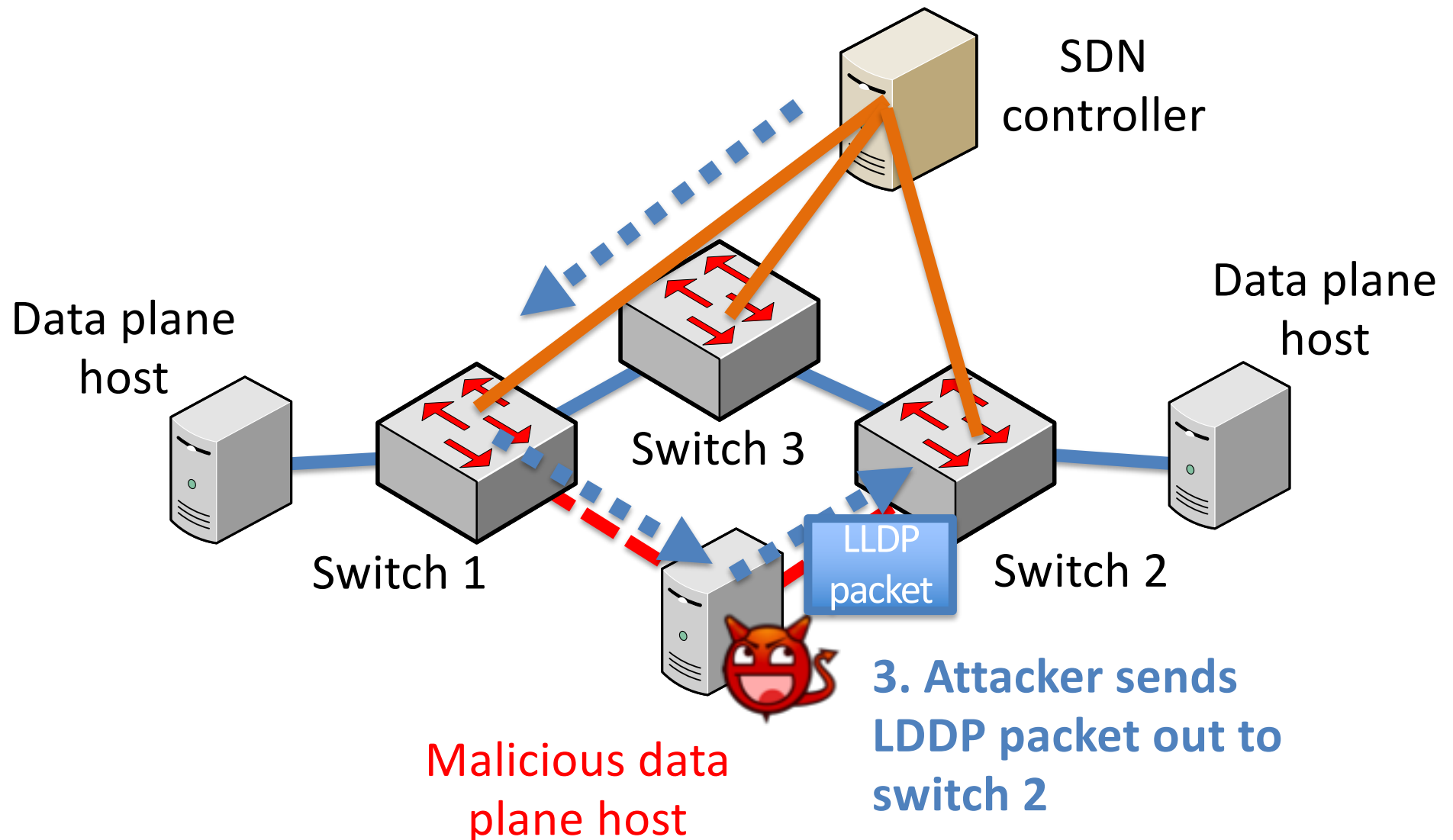
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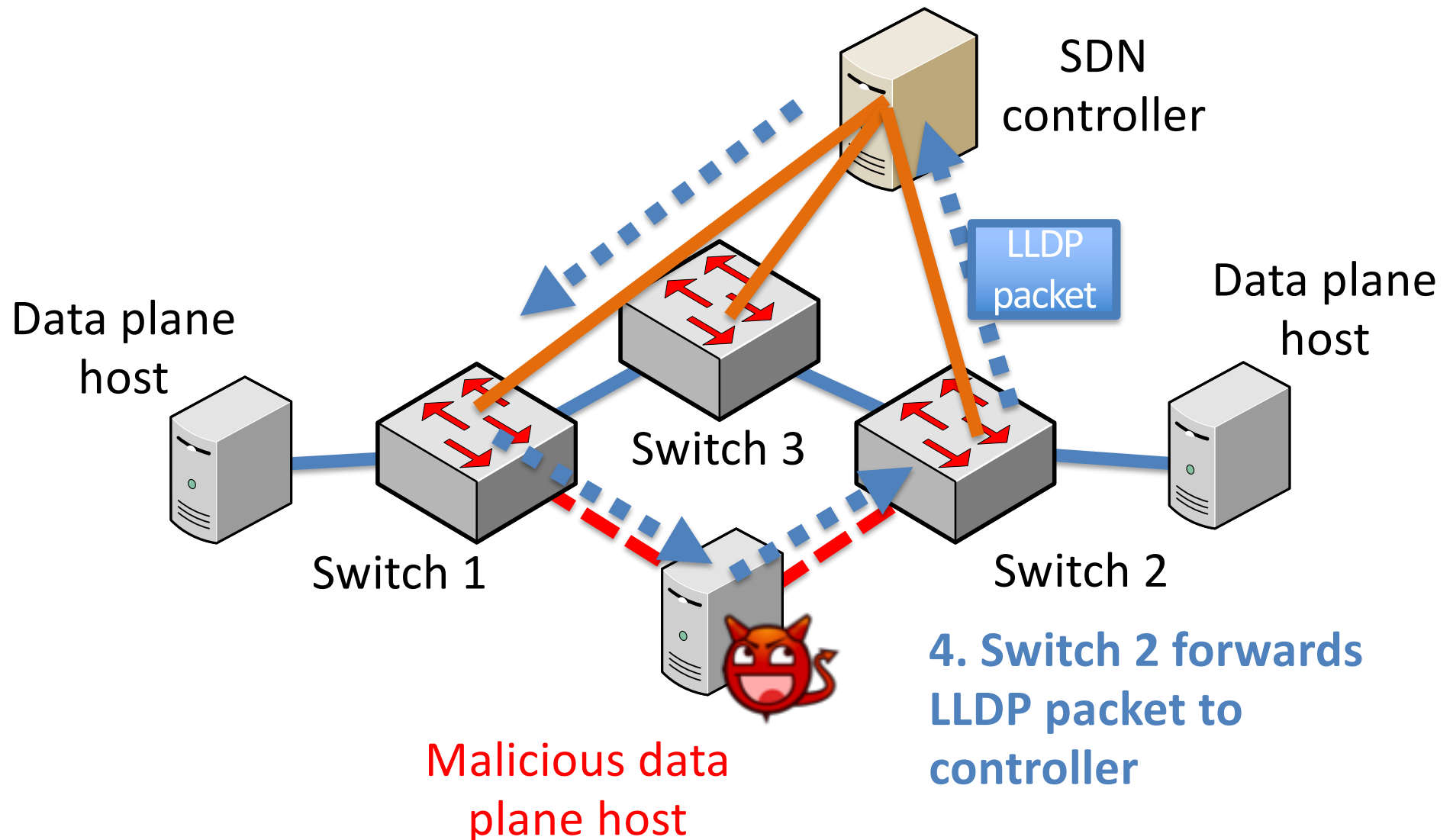
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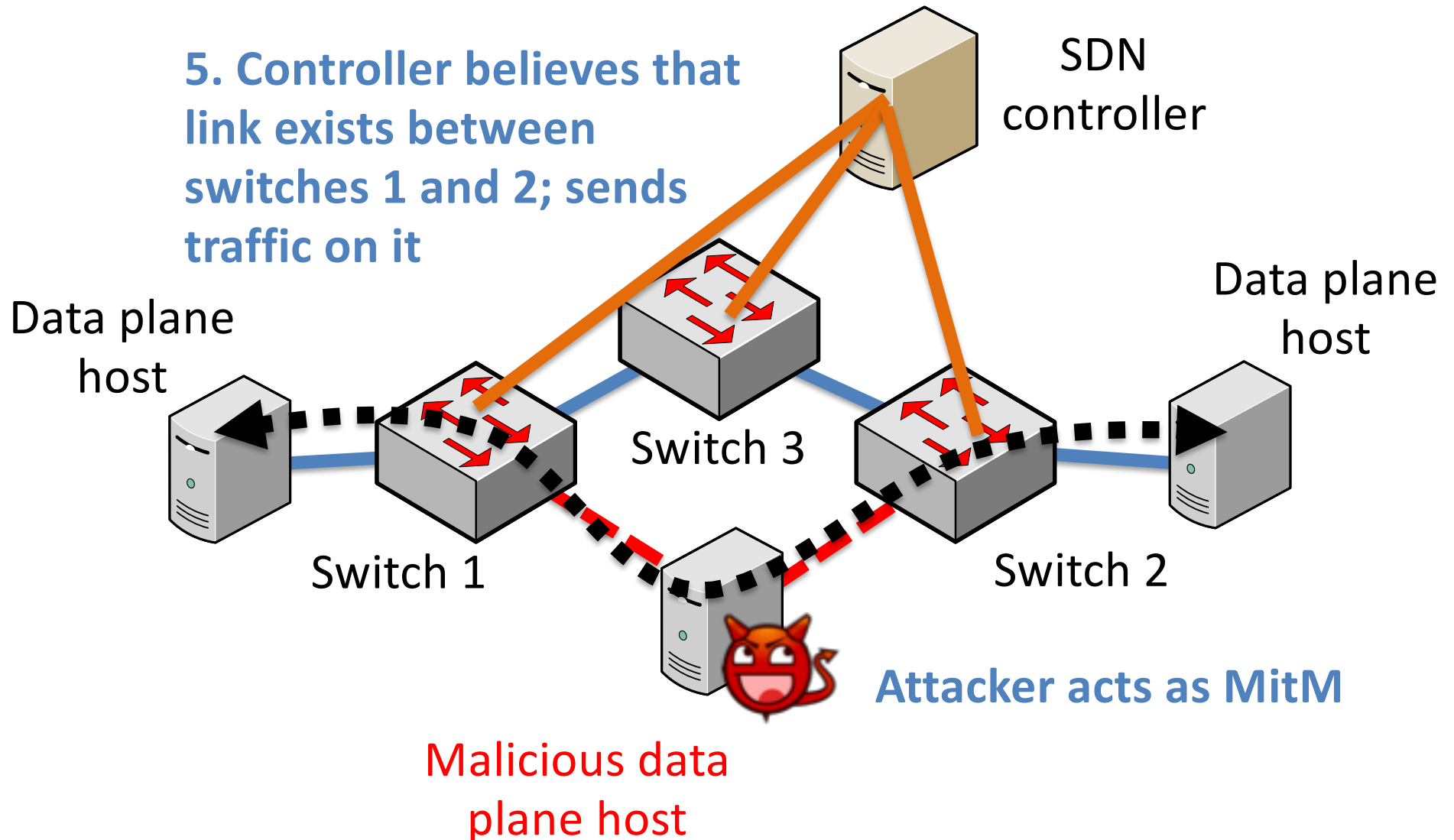
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# SDN Attacks: LLDP Spoofing

“Poisoning Network Visibility in Software-Defined Networks: New Attacks and Countermeasures” [NDSS ‘15]

- In SDN: the **link fabrication attack**
- Why this attack works:
  - SDN controllers maintain a topology graph (links, switches, and hosts) of the current network “state” so network applications can query it
  - LLDP packets must be sent through the data plane
  - Integrity of LLDP packets are (generally) not checked
    - Spoof, replay, or tunnel LLDP packets

# SDN Defenses: LLDP Spoofing

“Poisoning Network Visibility in Software-Defined Networks: New Attacks and Countermeasures” [NDSS ‘15]

- Spoofing: Signed LLDP messages
  - Uses Hash-based Message Authentication Code (HMAC) signed by the SDN controller
  - Provides integrity and authenticity guarantees that the LLDP packet came from the SDN controller
- Tunneling: Verify if port is HOST or SWITCH
  - Host and switch ports are mutually exclusive
  - Hosts will generate their own traffic (e.g., ARP, DNS), so these are not SWITCH ports

# SDN Security Research

- Active area of research!
- SDN as network operating system
  - Complexity of SDN framework creates new security challenges
  - What security concerns from OS design can we incorporate into SDN controller design?
- SDN application trust
  - SDN apps now have “app stores”
  - What assumptions should we make about trustworthiness of SDN apps?

# The End

- Questions?
- Announcements:
  - Networking Project, Checkpoint 1 due 6pm
  - Adam returning on Friday (lecture on TLS)