

- (A) $\pi_{Eid, Pid}(\text{Works}) / \pi_{Pid}(\sigma_{\text{Name} = 'M'}(\text{Project}))$
 (B) $\pi_{Eid}(\text{Works}) - \pi_{Eid}[\pi_{Eid}(\text{Works}) \times \pi_{Pid}(\sigma_{\text{Name} = 'M'}(\text{Project}))] - \pi_{Eid, Pid}(\text{Works})$
 (C) $\pi_{Eid}(\text{Works}) - \pi_{Eid}[\pi_{Eid}(\text{Works}) \times \pi_{Pid}(\sigma_{\text{Name} \neq 'M'}(\text{Project}))] - \pi_{Eid, Pid}(\text{Works})$
 (D) None of the above

Q8 Consider the two relations R_1 and R_2 such that they have no attributes in common.

$$S_1: R_1 \bowtie R_2 = R_1 \times R_2$$

$$S_2: R_1 \bowtie R_2 = \emptyset$$

Which of the following is correct?

- (A) S_1 only
 (B) S_2 only
 (C) Both S_1 and S_2 only
 (D) Neither S_1 nor S_2

Q9 Suppose that two relations $R(A, B)$ and $S(A, B)$ have exactly the same schema. Which of the following is/are always true?

- (A) $R \cap S = R - (R - S)$
 (B) $R \cap S = R - (S - R)$
 (C) $R \cap S = R \bowtie S$
 (D) $R \cap S = \sigma_{R.A=S.A \wedge R.B=S.B}(R \times S)$

Q10 Consider two relations $R(A,B)$ and $S(B,C)$ and the

following relational algebra expression $S: \pi_{R.A, S.B}(\sigma_{R.B=S.B}(R \times S))$.

Which of the following relational algebra expression(s) are guaranteed to produce same result as S ?

- (A) $\pi_{A,B}(R \bowtie S)$
 (B) $R \bowtie \pi_B(S)$
 (C) $R \cap (\pi_A(R) \times \pi_B(S))$
 (D) $\pi_{A, R.B}(R \times S)$

Q11 Consider the following relational schemas

$\text{Supplier}(\underline{\text{Sid}}, \text{Sname})$, $\text{Parts}(\underline{\text{Pid}}, \text{Pname})$,
 $\text{Catalog}(\underline{\text{Sid}}, \underline{\text{Pid}}, \text{Cost})$ and following two queries.

$$Q1: \pi_{\text{Sid}}[\text{Catalog} - [\pi_{C1.Sid, C1.Pid, C1.cost}(\sigma_{C1.cost < C2.cost}(\rho_{C1}(\text{Catalog}) \times \rho_{C2}(\text{Catalog})))]]$$

$$Q2: \pi_{\text{Sid}}[\text{Catalog} - \pi_{\text{Sid}}[\pi_{C1.Sid, C1.Pid, C1.cost}(\sigma_{C1.cost < C2.cost}(\rho_{C1}(\text{Catalog}) \times \rho_{C2}(\text{Catalog})))]]$$

Which of the following option is correct?

- (A) Both Q1 and Q2 always produces the same output.
 (B) Output produced by Q1 is always different from output produced by Q2.
 (C) Output produced by Q1 is subset of output produced by Q2.
 (D) Output produced by Q2 is subset of output produced by Q1.



GATE

Answer Key

Q1 (B)

Q2 (A)

Q3 (C)

Q4 (C)

Q5 (B)

Q6 (B)

Q7 (A, B)

Q8 (A)

Q9 (A, C)

Q10 (A, B, C)

Q11 (D)



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Hints & Solutions

Q1 Text Solution:

S_1 : Projection is not commutative.

S_2 : Selection is commutative.

Q2 Text Solution:

The expression $I \bowtie J \bowtie K$ will produce the following result:

| p | q | r | s |
|---|---|----|----|
| 0 | 1 | 2 | 3 |
| 4 | 5 | 2 | 3 |
| 4 | 5 | 6 | 7 |
| 4 | 5 | 10 | 11 |
| 4 | 5 | 10 | 3 |

Q3 Text Solution:

Sid who only enrolled for Papercode having description as "CS" = All enrolled Sid – All Sid who have enrolled in at least one Papercode with a description other than "CS"

Q4 Text Solution:

The employee ids who work in exactly one project = All employee ids who work in some project - All employee ids who work in at least two projects

Q5 Text Solution:

- I. Returns all values of B except the smallest one.
- II. Returns the smallest value of B.
- III. Returns the largest value of B.
- IV. Returns all values of B.

Q6 Text Solution:

The join condition ($Marks < M \wedge Gender = G$) filters out students who have lower marks than someone else of the same gender.
The set difference then removes all students who have marks lower than someone else in their gender.

The remaining sids are those who have the highest marks within their respective gender.

Q7 Text Solution:

Option A and option B are equivalent where both of them displays the Eids who work in every project with Name = 'M'.

Q8 Text Solution:

When no common attributes exist, a natural join behaves like Cartesian product.

Q9 Text Solution:

- A) This is a valid set identity. The intersection of two sets can be expressed as the elements of R after removing elements that are not in S.
- B) The expression $S - R$ gives elements in S but not in R. Subtracting these from R does not necessarily give $R \cap S$, so this identity does not always hold.
- C) Since R and S have the same schema, a natural join between them will return only common tuples, which is precisely $R \cap S$.
- D) The result will contain pairs of identical tuples, rather than a single instance of common tuples.

Q10 Text Solution:

$R \times S$ creates all possible pairs of tuples.
Projecting A, R.B means we include all possible combinations, not just valid join conditions.
Thus, option D is incorrect.

Q11 Text Solution:

Q1 removes specific (Sid, Pid, Cost) entries, meaning that a supplier could still remain if they have other entries in the catalog.
Q2 removes entire suppliers if they have at least one entry where they are not the cheapest.
Therefore, Q2 removes more suppliers than Q1, making Q2 a subset of Q1.


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