

CS & IT ENGINEERING



Operating System

Process Synchronization

Lecture -1

By- Vishvadeep Gothi sir



Recap of Previous Lecture



Topic

Multithreading

Topic

System Call: Fork()

Topics to be Covered



Topic

Synchronization

Topic

Race Condition

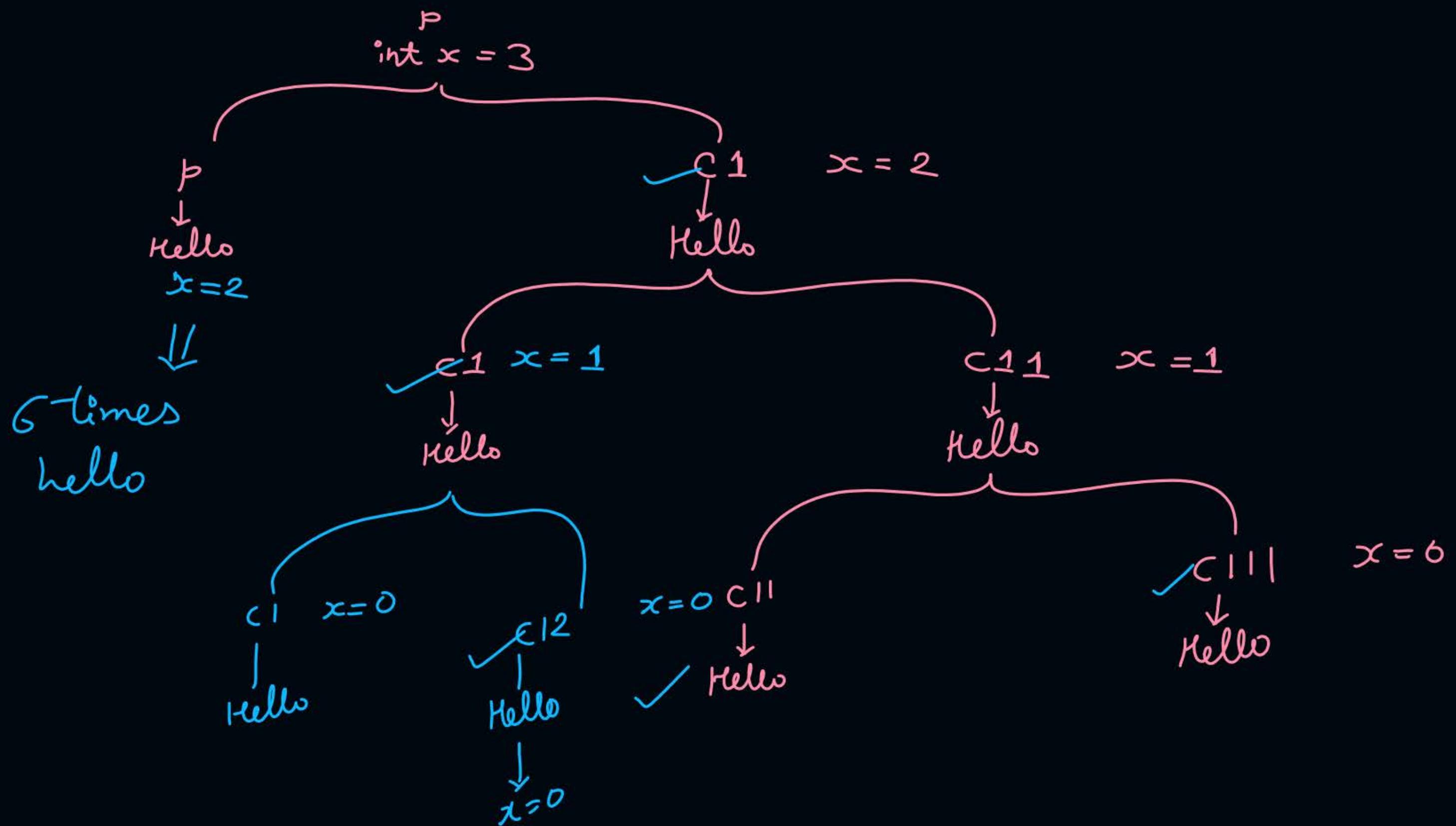
Topic

Critical Section

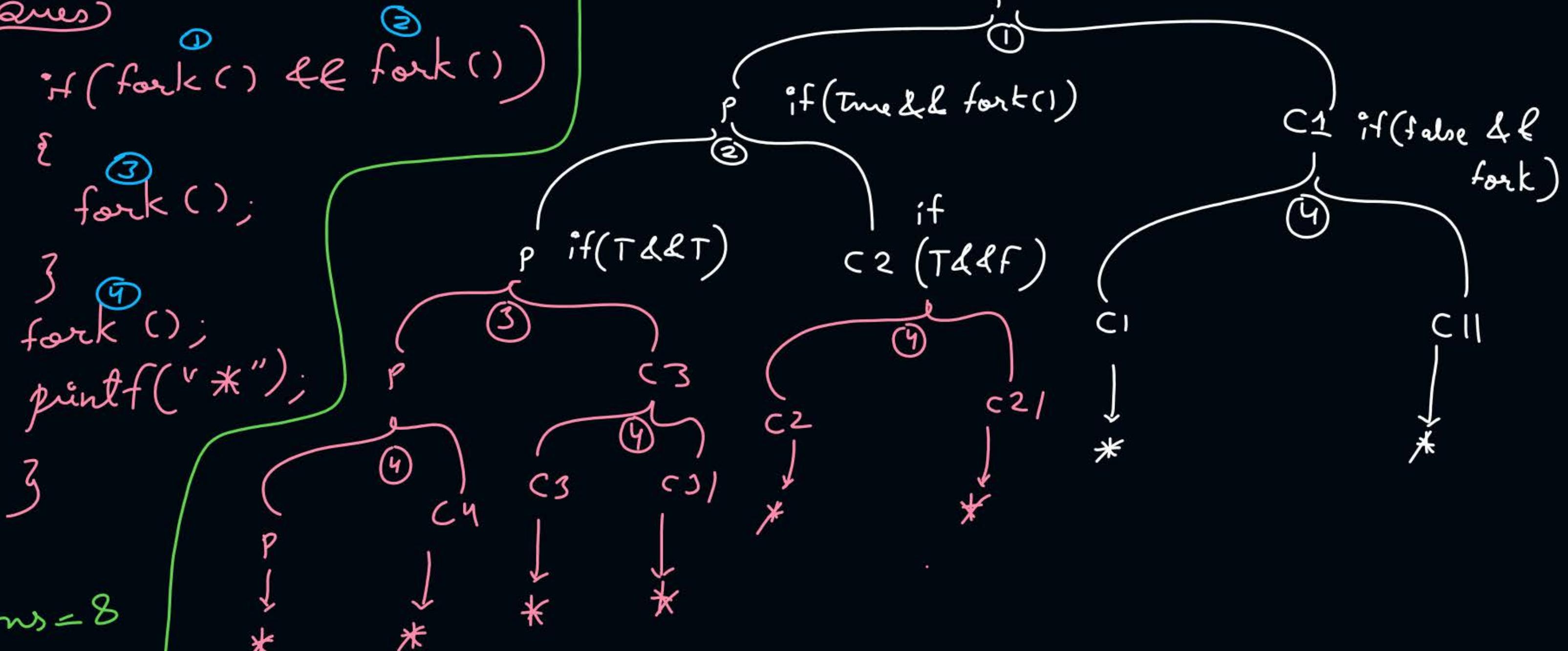
#Q. Consider the following code snippet using the fork() and wait() system calls. Assume that the code compiles and runs correctly, and that the system calls run successfully without any errors.

```
int x = 3;  
while(x > 0) {  
    fork();  
    printf("hello");  
    wait(NULL);  
    x--;  
}
```

The total number of times the printf statement is executed is 14?



Ques)



Ans = 8



Topic : Process Types

1. Independent
2. Cooperating/Coordinating/Communicating





Topic : Problems Without Synchronization



Problems without Synchronization:

- Inconsistency
- Loss of Data
- Deadlock

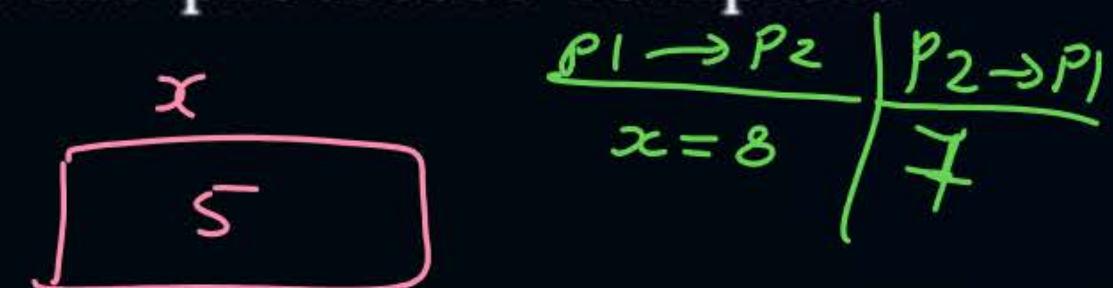


Topic : Race Condition



A race condition is an undesirable situation, it occurs when the final result of concurrent processes depends on the sequence in which the processes complete their execution.

due to lack of synchronization



P1
 $R1 \leftarrow 5$
 $\cancel{R1} \leftarrow R1 + 2$
 $x \leftarrow R1$

P2
 $R2 \leftarrow x$
 $R2^8 \leftarrow R3 + 3$
 $x \leftarrow R3$

[NAT]

P
W

#Q X = 10

P1
X = X / 2
5

P2
X = X + 4
14

Case 1 :- P1 then P2 9

Case 2 :- P2 then P1 7

Case 3 :- both read x concurrently then P2 writes last 14

Case 4 :- ———(1) ————— P1 ———
5

How many different values of X are possible after both processes finish executing? Ans = 4 (5, 7, 9, 14)

#Q. The following pair of processes share a common variable X.

Process A	Process B
int Y;	int Z;
$Y = \cancel{X}^{\cancel{6}} * 2;$	$Z = \cancel{X}^{\cancel{4}} + 1;$
$X = Y;$	$X = Z;$

9, 10, 8, 5

X is set to 4 before either process begins execution. As usual, statements within a process are executed sequentially, but statements in process A may execute in any order with respect to statements in process B.
How many different values of X are possible after both processes finish executing?

[NAT]

The following two functions $P1$ and $P2$ that share a variable B with an initial value of 2 execute concurrently.

```
P1 () {  
    C = B - 1;  
    B = 2 * C;  
}
```

```
P2 () {  
    D = 2 * B;  
    B = D - 1;  
}
```

3, 4, 2

The number of distinct values that B can possibly take after the execution is _____.

3

[NAT]

Consider three concurrent processes P1, P2 and P3 as shown below, which access a shared variable D that has been initialized to 100.

$$D = +20 +30 -50$$

P1	P2	P3
:	:	:
100	100	100
<u>D = D + 20</u>	D = D - 50	<u>D = D + 10</u>
:	50	130
:	:	120
:	:	:

$$\begin{aligned} Y &= 130 \\ X &= 50 \\ \hline \text{Ans} &= 80 \end{aligned}$$

The processes are executed on a uniprocessor system running a time-shared operating system. If the minimum and maximum possible values of D after the three processes have completed execution are X and Y respectively, then the value of Y - X is 80.

ques) In prev. quest', no. of distinct possible values D can take ?

Ans = 7

Soln

$$100 + 20 - 50 + 10$$

all 3 $\Rightarrow 80$

any 2 $\Rightarrow 70, 130, 60$

any 1 $\Rightarrow 120, 50, 110$

7 distinct values

ques)

$$x = 20$$

p_1

$$x = x + 1$$

p_2

$$x = x + 2$$

p_3

$$x = x + 4$$

p_4

$$x = x + 8$$

21, 22, 24, 28
23, 25, 29, 26, 30, 34
27, 32, 33, 31, 35

}

15 distinct value

ques) In prev. quest'n if value x can be operated only by one process at a time, then distinct possible values of $x _ 1$?

value $\Rightarrow 35$

↓
mutual exclusion

ques) $x = 20$

$$\frac{P_1}{x = x - 1}$$

$$\frac{P_2}{x = x + 2}$$

$$\frac{P_3}{x = x - 2}$$

$$\frac{P_4}{x = x + 3}$$

1. If no mutual exclusion then

$\max \frac{25}{17} ?$
 $\min \frac{17}{25} ?$
no. of distinct values 9 ?

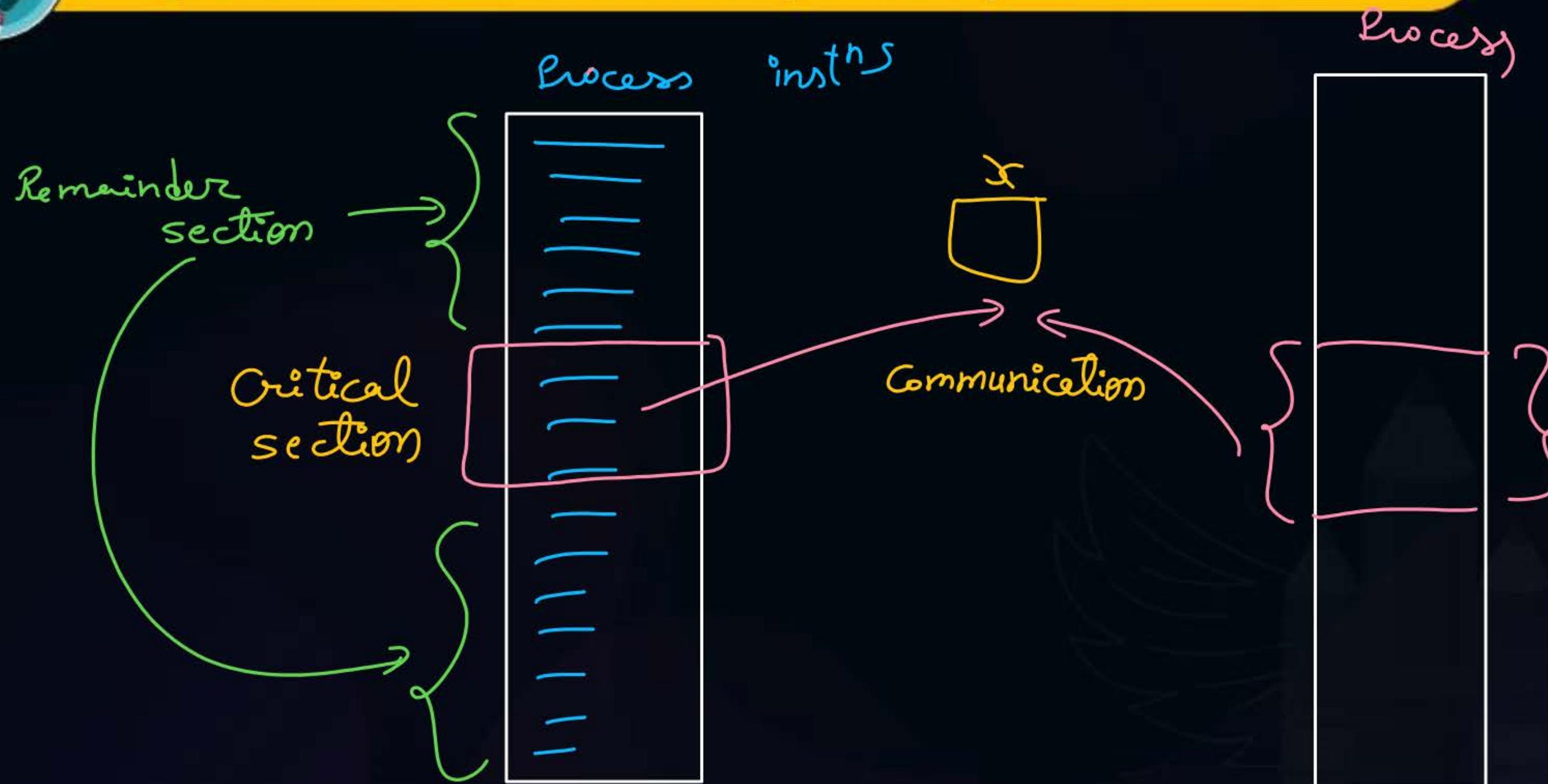
2. If mutual exclusion then $\Rightarrow \max_{\min} \{22, 1\}$
Distinct $\Rightarrow 1$

$$20 - 1 + 2 - 2 + 3$$

$$\begin{array}{r} 19, 22, 18, \\ 23, 21, 17, 20, 25, \\ 24, \end{array}$$



Topic : Entire Process Requires Synchronization?

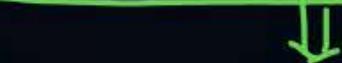




Topic : Critical Section



The critical section is a code segment where the shared variables can be accessed.



Communication happens .



Topic : Solution of Critical Section Problem



Requirements of Critical Section problem solution:

1. Mutual Exclusion
2. Progress
3. Bounded Waiting



Topic : Solution of Critical Section Problem



Mutual Exclusion:

If one process is executing the critical section, then other process is not allowed to enter into critical section.



Topic : Solution of Critical Section Problem



Progress:

If no any process is in critical section and any process wants to enter into critical section, then the process must be allowed.



Topic : Solution of Critical Section Problem



Bounded Waiting:

If a process p₁ is executing in critical section and other process p₂ is waiting for critical section, then the waiting time of p₂ must be bounded. Which means p₁ must not enter in to critical section again and again by keeping p₂ in waiting for long.



2 mins Summary

Topic

Synchronization

Topic

Race Condition

Topic

Critical Section





Happy Learning

THANK - YOU