



CS & DA



Database Management System

File Organisation & Indexing

DPP 01 (Discussion Notes)



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[MCQ]

#Q.
the

Assume a relational database system that holds relation: C(colleges) with following characteristics

- Records are stored as fixed length, fixed format records, length is 256 bytes.
- There are 16384 records.
- Records contains key attribute CollegeNumber (C.N), length 22 bytes and other fields.
- Unspanned organization is used to store the information or record.

$$\# \text{ blocks of relation } C = \frac{16384}{16} = 1024$$

Let's suppose we want to build a sparse primary index on C.N then how many numbers of 4096-byte blocks are needed to store the primary index when block pointer size is 10 bytes _____?

A

7

$$\begin{aligned} \text{Size of index entry} &= 22 + 10 = 32 \\ \# \text{index records per block} &= \frac{4096}{32} \\ &= 128 \end{aligned}$$

C

9

B

8

$$\# \text{ records in a block} = \frac{4096}{256}$$

D

10

$$\begin{aligned} \# \text{ blocks to store PI} &= \frac{1024}{128} = 8 \\ &\checkmark 16 \end{aligned}$$

#Q. Assume a relational database system that holds relation: Product (P) with the following characteristics

- Records are stored as fixed length, fixed format records, with the length of 256 bytes.
- There are 262144 records.
- Records contain attribute P.I (The identifier of the product involved), with the length 24 bytes, and an attribute P.C (the cost of product), with the length 32 bytes and other fields.
- Unspanned organization is used to store the record.

Assume that we want to build a dense secondary index on P.C, then how many numbers of 4096-byte blocks needed to store the dense secondary index. When record pointer size is 32 bytes? 4096.

$$\# \text{ blocks} = \frac{262144}{64}$$

$$\# \text{ entries in index block} = \# \text{ records}$$

$$= 262144$$

[MCQ]

#Q. Consider a SQL statement (SELECT P₁, P₂, P₃ from Q WHERE P₂ = 'Pavan') is frequently executed, which column(s) should be considered for indexing based only on the statement itself?

- A** P₁ only
- B** P₂ only
- C** P₃ only
- D** P₁, P₂ and P₃

[MCQ]

#Q. Consider the following specification of system-Disk block size = 2048 bytes

Block pointer size = 16 bytes

Record pointer size = 20 bytes long file contains 30,000 records.

Each record of the file has the following fields:

An extra/additional byte is used per record to represent end of the record.

What is the block factor of the database file assuming unspanned file organization?

A 16

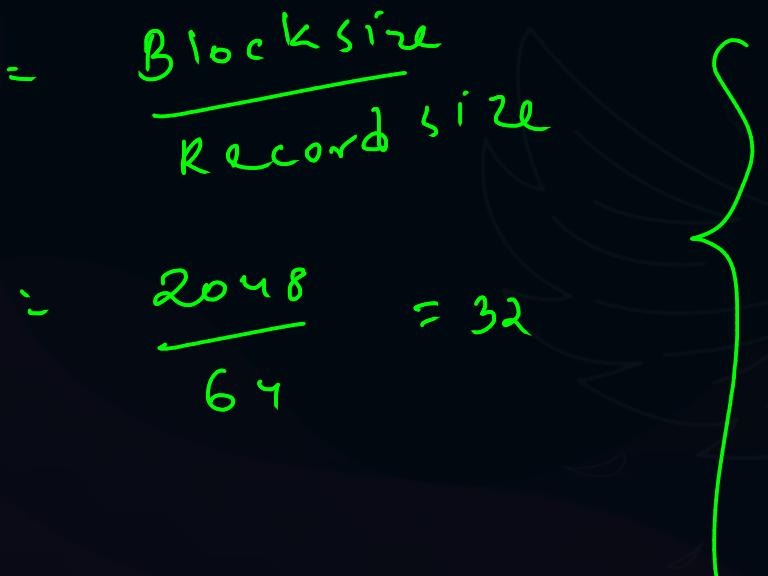
$$\text{Block factor} = \frac{\text{Block size}}{\text{Record size}}$$

$$= \frac{2048}{64} = 32$$

B 32

C 48

D 64



Fields	Size (in Bytes)
EmpName	5
EmpNum	10
DeptNum	9
Addr	20
PhNum	9
DOB	1
Sex	1
Job	3
Sal	5

$\{ \text{size} = 63$
 $+ 1 = 64$

#Q. Which one of the following statements is/are True regarding indexing?

- A A database file can contain multiple clustered indexes.
- B A database file can consist of at most one clustered index with multiple secondary indexes.
- C A database file can consist of multiple primary indexes.
- D A database file can consist of both primary and clustered index.

[NAT] # block accesses = search in index file + 1st-block + 2nd block
 $= \lceil \log_2 74 \rceil + 1 + 1 = 7 + 1 + 1 = 9$

#Q. Consider a database of fixed-length records stored as an ordered file. The database has 25,000 records with each record being 100 bytes, of which the non-key attribute on which clustering index is formed occupies 10 bytes. The data file is completely block aligned.

Suppose, block size, of the file system is 512 bytes and a pointer to the block occupy 5 bytes. You may assume that a binary search on an index file of b block may take $\lceil \log_2 b \rceil$ accesses in worst case.

Given that a cluster consumes 2 blocks, the number of block accesses required to identify the desired data in the worst case is ____.

Block factor of DB file = $\left\lceil \frac{512}{100} \right\rceil = 5 \text{ records/block}$

blocks to store 25000 records = $\frac{25000}{5} = 5000 \text{ blocks}$

entries in index file = $\frac{5000}{2} = 2500$

Block factor of index file = $\left\lceil \frac{512}{15} \right\rceil = 34$

blocks in index file = $\left\lceil \frac{2500}{34} \right\rceil = 74$

[MCQ]

#Q. Consider the following statements-

- S₁:** If the records of a relation X are physically ordered over a non-key field P and an index is build over the key-field of relation X, then the index is necessarily a secondary index over key attribute.
- S₂:** More than one secondary indexes are possible.

Which of the given statement(s) is/are CORRECT?

A

S₁ only

B

S₂ only

C

Both S₁ and S₂

D

Neither S₁ nor S₂

#Q. The order of a leaf node in a B+ tree is the maximum number of (value, data record pointer) pairs it can hold. Given that the block size is 1K bytes, data record pointer is 8 bytes long, the value field is 10 bytes long and a block pointer is 6 bytes, then what is the order of the leaf node ?

A

53

C

55

B

54

D

56

$$B + P(K + R) \leq 10^24$$

$$\Rightarrow 6 + P(10 + 8) \leq 1024$$

$$\therefore =$$

$$P = \left\lfloor \frac{10^{18}}{18} \right\rfloor = 56$$

#Q. (The order of an internal node in B+ tree index is the maximum number of children it can have.) Assume that a child pointer takes 6 bytes, the search field value takes 34 bytes and the blocks size is 2048 bytes. The order of the internal node is 52.

$$(p-1)34 + 6p \leq 2048$$

$$p \leq \frac{2082}{40}$$

$$p = \lfloor 52.05 \rfloor = 52$$

#Q. Given a block can hold either 3 records or 10 key pointers. A database contains P records, then how many blocks do we need to hold the data file and the dense index ?

A $\frac{P}{30}$

C $\frac{13P}{30}$

B $\frac{P}{3}$

D $\frac{P}{10}$

blocks storing
records = $\frac{P}{3}$

blocks storing
keys = $\frac{P}{10}$

Total no of blocks = $\frac{P}{3} + \frac{P}{10}$

= $\frac{13P}{30}$

#Q. Assume a disk with block size $B = 1024$ Bytes, A block pointer is $PB = 12$ bytes long and a record pointer is $PR = 18$ bytes long. A file has 1,00,000 patients records of size 100 bytes. Suppose the file is ordered by the key field PID and we want to construct a secondary (dense) index on non-key field Dept ID (14 bytes), then minimum of how many blocks are required to store index file assuming an unspanned organization?

$$BF : \frac{1024}{100} = 10$$

A

*Index record
size = 14 + 18
= 32 B*

C

B

3100

$$\# \text{ blocks} = \frac{100000}{10} = 10,000$$

D

None of the above

$$BF_{ind} = \frac{1024}{32} = 32$$

$$1st \text{ level index blocks} = \frac{1000000}{32} = 3125$$

#Q. The order of a node in B tree is the maximum number of block pointers it can hold. Given that the block size is 2K bytes, data record pointer is 8 bytes long, the search key is 9 bytes long and a block pointer is 5 bytes long. The best possible order of B tree node is .

$$\text{Block size} \geq 5P + (P-1)(8+9)$$

$$\Rightarrow 2048 \geq 5P + (P-1) \cdot 17$$

$$P = \left\lfloor \frac{2048}{22} \right\rfloor = 93$$

#Q. Consider the keys (1- 5000) are going to be inserted into a B⁺ tree. Assume, all the order are available before insertion. The orders P for B+ tree node is defined as-

2 to P pointer for root

$\left\lceil \frac{P}{2} \right\rceil$ to P pointer for another node.

The maximum possible levels in a B⁺ tree index for P = 9 is 6.

(Assume that level of the root node is 1).

$$\left\lceil \frac{2}{2} \right\rceil = 1 \quad \checkmark$$

$$\left\lceil \frac{1}{2} \right\rceil = 1 \quad \checkmark$$

$$\left\lceil \frac{10}{5} \right\rceil = 2 \quad \checkmark$$

$$\left\lceil \frac{50}{5} \right\rceil = 10 \quad \checkmark$$

$$\left\lceil \frac{250}{5} \right\rceil = 50 \quad \checkmark$$

$$\left\lceil \frac{1250}{5} \right\rceil = 250 \quad \checkmark$$

$$\text{last level} = \left\lceil \frac{5000}{4} \right\rceil = 1250 \quad \checkmark$$

#Q. Consider the following statements:

S₁: In a B+ tree, data pointers are stored only at the leaf nodes of the tree.

S₂: The leaf node has an entry for every value of the search field, along with the data pointer to the record.

Choose the correct statements.

**A**

Only S₁ is true

**C**

Both S₁ and S₂ are true

**B**

Only S₂ is true

**D**

Neither S₁ nor S₂ is true

#Q. Which of the following is/are true reading B+ tree ?

- A Records can be fetched in equal number of disk access.
- B Height of the tree remains balanced and less as compared to B tree.
- C Keys are used for indexing
- D Faster search queries as the data is stored only on the leaf nodes.

#Q. Consider the following specification of system with disk block size 2048 bytes, block pointer size 14 bytes, record pointer size 18 bytes long and file size 60,000 records. Each record of file is 256 bytes long and record of the size is sorted on the key field. If the primary index (sparse) is built on the kye field (ESN) which is 18 bytes long. What is the Index blocking factors (That is number of indexes per block) .

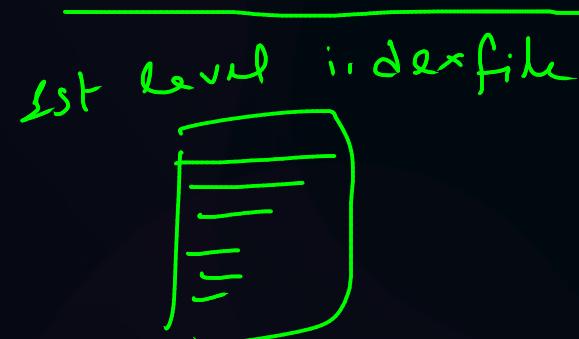
Assuming unspanned file organization 64.

$$\# \text{ indexes per block} = \frac{2048}{14 + 18} = 64$$

#Q. Data For This & Next Question:

Consider a disk blocking size $B = 1024$ bytes. A block pointer is $PB = 12$ bytes long and a record pointer is $PR = 7$ bytes long. A file has $r = 60,000$ patient records of fixed length. The size of record is 230 bytes. Suppose the file is not ordered by the key field PSN (18 bytes) and we want to construct a secondary index on PSN.

The number of first level index entries is 60,000.



[MCQ]



Index Record

$$\text{size} = 18 + 14 = 32 \text{ B}$$

#Q. The number of first level index block is _____?

A 1800

C 1850

B 1825

D 1857
1875

$$BF_{\text{ind}} = \left\lceil \frac{1024}{32} \right\rceil = 32$$

$$\# \text{ 1st level index block} = \frac{60000}{32}$$

$$= 1875$$

#Q. Consider an unordered file of 10^6 records with records size of 200 bytes stored on blocks of 8KB with a spanned records organization. We will assume that no system related information is stored within a block, then how many blocks would it be need to store this file?

A 24400

C 24410

B 24405

D 24415

$$\# \text{records in a block} = \frac{8KB}{200} = 40.96$$

$$\# \text{blocks} = \frac{10^6}{40.96} = 24415$$

#Q. Consider the following statements:

- X S_1 : for any given data file, it is possible to create two different sparse first level indexes on various keys.
- ~~S_2 : for any given data file, it is possible to create two different denes first level indexes on various keys.~~ Select the correct statements.

A

Only S_1 correct

B

Only S_2 correct

C

Both S_1 and S_2 are correct

D

Neither is S_1 and S_2 is correct



THANK - YOU