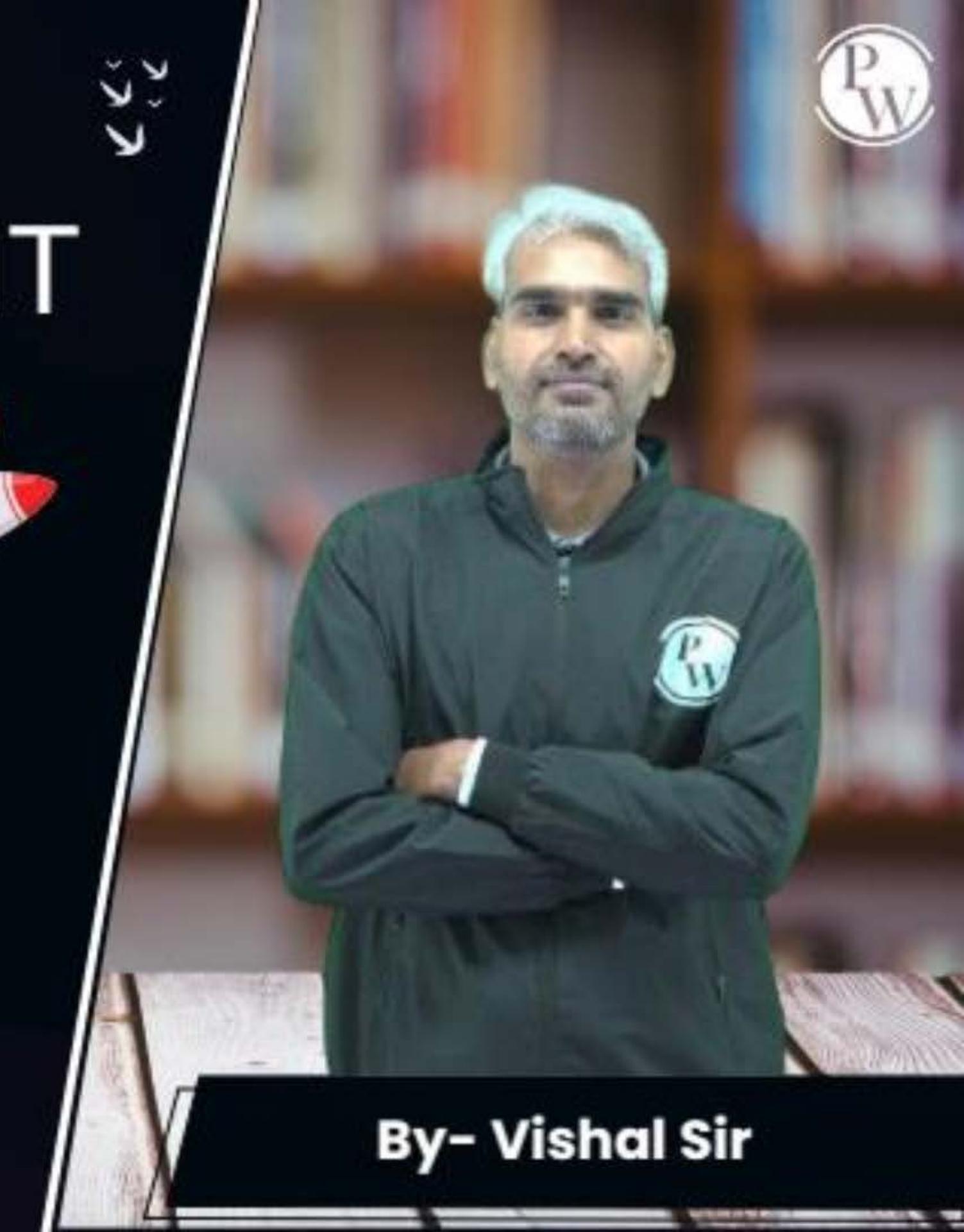


Computer Science & IT

Database Management System

Query Languages

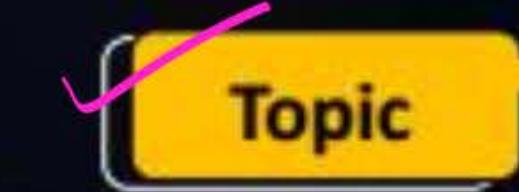
Lecture No. 07



By- Vishal Sir

Recap of Previous Lecture



- Topic SQL clauses
- Topic Nested query (sub-query)
- Topic Operators w.r.t. Nested query
- Topic Order of execution of nested query

Topics to be Covered



Topic

Execution of correlated sub-query

Topic

AS clause

Topic

WITH clause

Nested Query

✓ Independent Nested query
(Independent Sub-query)

* If inner query can be executed independently, then it is called independent sub-query

✓ Correlated Nested query
(Correlated Sub-query)

→ When execution of inner query requires value of an attribute from the relation specified in Outer query, then it is called Correlated Sub-query

Nested Query

Independent Nested query (Independent Sub-query)

e.g. `Select Sid
from Student
Where (Marks) = (Select Max(Marks)
From Student)`

Main Query **Inner Query**

This query can
be executed
independently
∴ Independent
Sub-query

Correlated Nested query (Correlated Sub-query)

e.g. `Select R.A
From R
Where ... (Select S.C
From S
Where (R.A=S.B))`

Main Query **Inner Query**

Value of Attribute 'A' is required
from relation R (i.e., from the relation
specified in outer query)
∴ Correlated Sub-query

(1) Order of execution w.r.t. Independent Sub-query

- ④ Inner query will be executed first and it will produce its output,
- ④ then outer query will use the output produced by inner query for its execution.

* Retrieve Sids of all the Students

Who scored maximum Marks.

Select Sid
From Student

Where (Marks = (Select Max(Marks)
From Student))

O/P of this query = 60

And then

Our query will compare
the marks of each tuple
with "60"

Op :- Sid
S3
S4

Student

Sid	Sname	Marks	Branch
X S1	A	40	CS
X S2	A	20	IT
✓ S3	B	60	CS
✓ S4	A	60	EC
X S5	C	40	IT
X S6	C	NULL	EC

Correlated Sub-query :-

operand of this operator is also the o/p of inner query.

Main query
Select *
From R

Where operator

This operator
will return
True or False

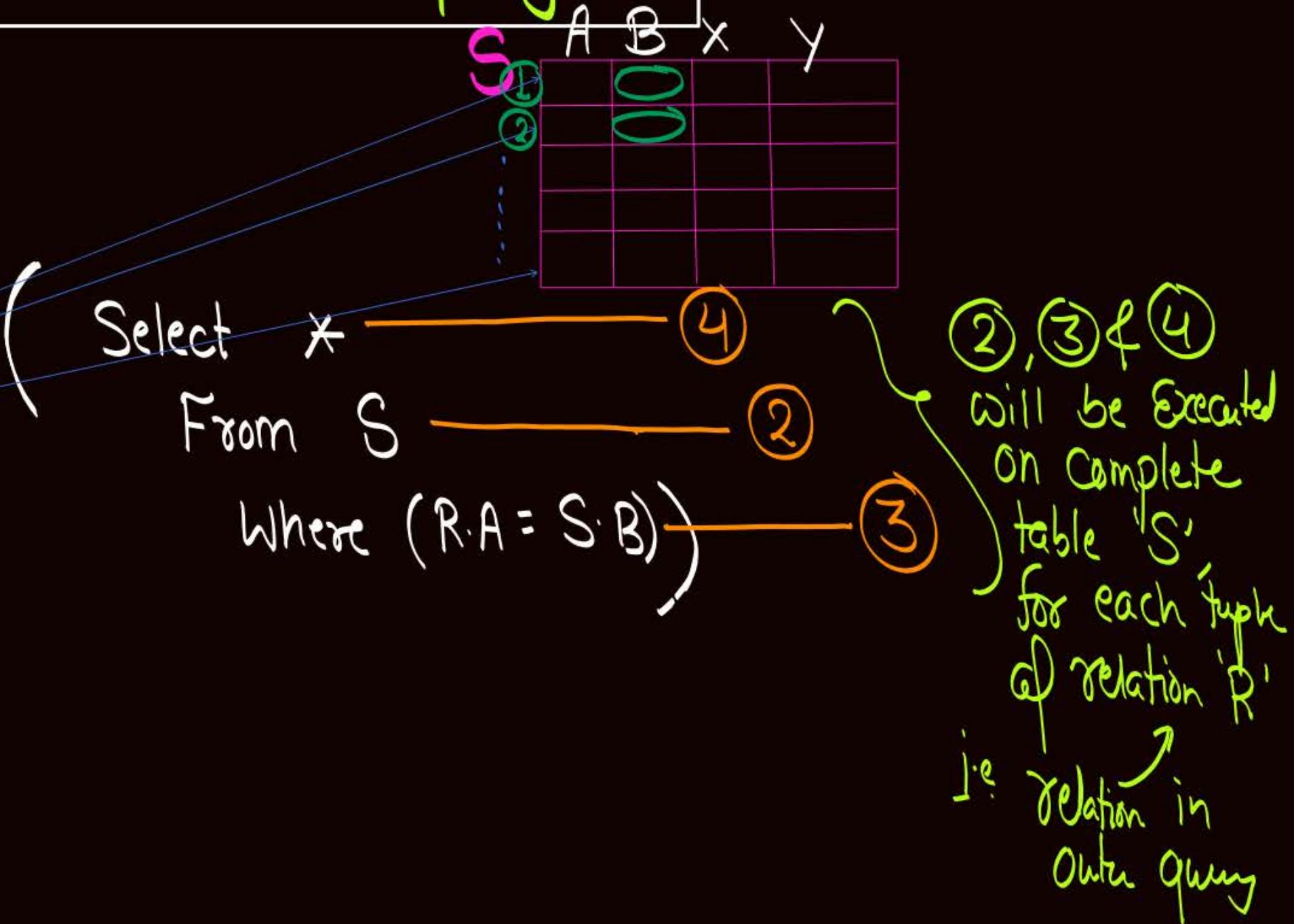
Inner query
Select *
From S
Where (R.A = S.B)

- If operator returns true, then "Where" Cond is true
- If operator returns false, then 'Where' Cond is false

Order of Execution w.r.t. Correlated Subquery :-

⑥ — Select *
 ① — From R
 ⑤ — Where operator

R	A	P	Q
1	0		
2			
3			
4			
5			
6			





Topic : Operators

* IN, ANY, ALL and EXISTS are the operators that are generally used with Sub-query Concept

IN operator

Operator "IN" is used to check whether the concerned tuple is a member of a set of tuples produced by the inner query or not?

Complement of 'IN' is 'NOT IN'

Check whether

X

IN

{ 2, 5, 8, 9, 12, 13 }

Our concerned
tuple from
Outer table

Set of tuples
Produced by
inner query

- for eg.

Select Eid
From EMP

(i) If $X = 8$, then operator 'IN' will return true

(ii) If $X = 10$, then operator 'IN' will return false

Check whether

X

IN

$\{(1, A), (3, D), (4, D), (7, F)\}$

"tuple from
Outer table"

for eg

Select Sid, Sname
From Students

(i) If $X = (4, D)$, then IN returns True

(ii) If $X = (F, 7)$, then IN returns false.

$\sqsubset \neq (7, F)$

Note :-

If set of tuples is empty {i.e., O/P produced by inner query is Empty}, then operator 'IN' will always return False

and hence operator 'NOT IN' will always return true if inner query result is Empty.

Q: Consider the following relational scheme

Supplier (Sid, Sname, Rating)

Parts (Pid, Pname, Color)

Catalog (Sid, Pid, Cost)

* Retrieve Sid of suppliers who supplied some Red

Color Parts

Parts	
P1	Red
P2	Green
P3	Red

Catalog	
S1	P1
S1	P2
S2	P2
S3	P2
S3	P3

R.A:-

$\pi_{Catalog.Sid} \left(\begin{array}{l} \sigma_{Catalog.Pid = Parts.Pid} (Catalog \times Parts) \\ \wedge \\ Parts.Color = 'Red' \end{array} \right)$

SQL: (Without IN)

Select C.Sid

Catalog
renamed
as C

From Catalog C , Parts P
Where (C.Pid = P.Pid AND P.Color = 'Red')

SQL: (Using IN)

Select Sid From Catalog

Where Pid IN (Select Pid From Parts)
Where Color = 'Red'

Parts
renamed as P

'ANY' operator

Operator 'ANY' is used along with Comparison operators

<, ≤, >, ≥, =, <>

↑
Not equal

Check whether $x < \text{ANY}$

$x = 10$	x	x	x	x	✓
$x = 20$	x	x	x	x	x

{ 2, 5, 8, 9, 12, 14 }

If $x = 10$, then "ANY" will return True

If $x = 20$, then "ANY" will return False

'ANY' operator

Operator 'ANY' is used along with Comparison Operators
↓
<, ≤, >, ≥, =, <>

Note: Operator 'ANY' will return True, if and only if
at least one tuple in the set of tuples satisfy the
Comparison Condition

Note: If set of tuples is empty, then Operator 'ANY' will
always return False

'ALL' operator

Operator 'ALL' is also used along with Comparison Operators

<, <=, >, >=, =, <>

Check whether $x = 10$, $x < \text{ALL}$

$x = 10$, $x \in \{2, 5, 8, 9, 12, 14\}$

$x = 1$, $\underline{\quad \checkmark \quad \checkmark \quad \checkmark \quad \checkmark \quad \checkmark}$

If $x = 10$, then ALL returns False

If $x = 1$, then ALL returns True

'ALL' operator

Operator 'ALL' is also used along with Comparison Operators

<, <=, >, >=, =, <>

- * Operator 'ALL' can return False if and only if at least one tuple in the set of tuples fails the Comparison Condition

- * If set of tuples is Empty, then operator 'ALL' will always return True

EXISTS operator

- EXISTS returns True if and only if inner query result is not Empty.
 - If inner query result is not Empty, then EXISTS will return True &
 - If inner query result is Empty, then EXISTS will return False.

Complement of 'EXISTS' is 'NOT EXISTS'

Today's Topic

H.W.

#e.g.

```
SELECT C.sid
```

FROM Catalog C

WHERE EXISTS (SELECT *

FROM Parts P

WHERE (P.Pid=C.Pid AND P.color='RED')

Correlated
Sub-query

Parts	
Pid	Color
P ₁	Red
P ₂	Green
P ₃	Red

Catalog	
Sid	Pid
S ₁	P ₁
S ₁	P ₂
S ₂	P ₂
S ₃	P ₂
S ₃	P ₃

What output will be produced by the above query if it is executed on the given instances of Parts & Catalog table

H.W.

#e.g.

SELECT C.sid → **S₁** ✓

is selected

FROM Catalog C

WHERE EXISTS (SELECT *

FROM Parts P

WHERE (P.Pid=C.Pid AND P.color='RED')

Catalog (C)

Sid	Pid
S ₁	P ₁
S ₁	P ₂
S ₂	P ₂
S ₃	P ₂
S ₃	P ₃

Return True

∴ 1st tuple of catalog is selected

∴ inner query result is not empty
(P₁, Red)

Pid	Color
P ₁ ✓	Red
P ₂ ✗	Green
P ₃ ✗	Red

P

H.W.

#e.g.

SELECT C.sid → S₁ ✓
 is selected

FROM Catalog C

WHERE EXISTS (SELECT *

Catalog (C)

Sid	Pid
S ₁	P ₁
S ₁	P ₂
S ₂	P ₂
S ₃	P ₂
S ₃	P ₃

FROM Parts P

WHERE (P.Pid=C.Pid AND P.color='RED')

Return False

∴ 2nd tuple of catalog is not selected

w.r.t 2ⁿ tuple
 of catalog
 Q/P is empty

Pid	Color
P ₁	Red
P ₂	Green
P ₃	Red

P

H.W.

#e.g.

SELECT C.sid → **S₁** ✓
 is selected

FROM Catalog C

WHERE EXISTS (SELECT *

Catalog (C)

Sid	Pid
S ₁	P ₁
S ₁	P ₂
S ₂	P ₂
S ₃	P ₂
S ₃	P ₃

FROM Parts P

WHERE (P.Pid=C.Pid AND P.color='RED')

Return False

∴ 3rd tuple of catalog is not selected

w.r.t 3rd tuple
 of catalog
 Q/P is empty

Pid	Color
P ₁	Red
P ₂	Green
P ₃	Red

P

H.W.

#e.g.

SELECT C.sid → S₁ ✓
 is selected

FROM Catalog C

WHERE EXISTS (SELECT *

Catalog (C)

Sid	Pid
S ₁	P ₁
S ₁	P ₂
S ₂	P ₂
S ₃	P ₂
S ₃	P ₃

w.r.t 4th tuple
 of catalog
 Q.P is empty

FROM Parts P

WHERE (P.Pid=C.Pid AND P.color='RED')

Return False

∴ 4th tuple of catalog is not selected

Pid	Color
P ₁	Red
P ₂	Green
P ₃	Red

H.W.

#e.g.

SELECT C.sid
FROM Catalog C

 S_1 S_3

is selected

WHERE EXISTS (SELECT *

Catalog (C)

Sid	Pid
S_1	P_1
S_1	P_2
S_2	P_2
S_3	P_2
S_3	P_3

Return True

5th tuple of catalog is

Selected

FROM Parts P

WHERE (P.Pid=C.Pid AND P.color='RED')

o/p: (P_3 , Red)
not empty

Pid	Color
P_1	Red
P_2	Green
P_3	Red

~~H.W.~~

#e.g.

```
SELECT C.sid
FROM Catalog C
WHERE EXISTS ( SELECT *
```

Catalog

Sid	Pid
S ₁	P ₁
S ₁	P ₂
S ₂	P ₂
S ₃	P ₂
S ₃	P ₃

FROM Parts P

WHERE (P.Pid=C.Pid AND P.color='RED')

final op =

Sid
S ₁
S ₃

H.W.
#e.g.

```
SELECT C1.sid
FROM Catalog C1
WHERE NOT EXISTS (SELECT P.Pid
                   FROM Parts P
                   WHERE NOT EXISTS (SELECT C2.Sid
                                      FROM Catalog C2
                                      WHERE(C2.Pid=P.Pid AND C2.Sid=C1.Sid)))
```

What output is produced by above SQL query:

- A. Sids of suppliers who supplied some parts
- B. Sids of suppliers who supplied only proper subset of parts from all parts
- C. Sids of suppliers who supplied all parts
- D. Sids of suppliers who did not supply any part

#e.g.

~~HW:~~

SELECT C1.sid

FROM Catalog C1

WHERE NOT EXISTS (SELECT P.Pid

Catalog

Sid	Pid
S ₁	P ₁
S ₂	P ₂
S ₁	P ₂

Parts

Pid
P ₁
P ₂

FROM Parts P

WHERE NOT EXISTS (SELECT C2.Sid

FROM Catalog C2

WHERE(C2.Pid=P.Pid AND C2.Sid=C1.Sid)))

What output is produced by above SQL query:

- A. {S₁, S₂} Sids of suppliers who supplied some parts
- B. {S₂} Sids of suppliers who supplied only proper subset of parts from all parts
- C. {S₁} Sids of suppliers who supplied all parts
- D. ~~S~~ Sids of suppliers who did not supply any part

Select $C_1.Sid$
From Catalog C_1
Where Not Exists (

C_1	
Sid	Pid
S_1	P_1
S_2	P_2
S_1	P_2

Select $P.Pid$
From Parts P

P	
	Pid
	P_1
	P_2

Where Not Exists (Select $C_2.Sid$
From Catalog C_2
Where ($C_2.Pid = P.Pid$
AND
 $C_2.Sid = C_1.Sid$))

C_2	
Sid	Pid
S_1	P_1
S_2	P_2
S_1	P_2

- ∴ Not EXISTS will return false
- ∴ P_1 is not selected

W.r.t. 1st tuple of parts 'P' inner query result is not empty

Select $C_1.Sid$
From Catalog C_1
Where Not Exists

w.r.t 1st tuple of
Catalog (C_1) inner

query result is empty

∴ Not EXISTS will
return true.

& Sid of the 1st tuple
is selected

C_1	
Sid	Pid
S_1	P_1
S_2	P_2
S_1	P_2

Select
From Parts P

~~P_1 P_2~~

$P.Pid$

P	
Pid	
P_1	
P_2	

Where Not Exists

Select $C_2.Sid$
From Catalog C_2

Where $(C_2.Pid = P.Pid$

AND
 $C_2.Sid = C_1.Sid)$

C_2	
Sid	Pid
S_1	$P_1 \times$
S_2	$P_2 \times$
S_1	$P_2 \checkmark$

- ∴ Not EXISTS will return false
- ∴ P_2 is not selected

w.r.t 2nd tuple
of parts 'P' inner
query result is not empty

Select $C_1.Sid$
From Catalog C_1
Where Not Exists (

C_1	
Sid	Pid
S_1	P_1
S_2	P_2
S_1	P_2

P	
Pid	
P_1	
P_2	

Select $P.Pid$
From Parts P

Where Not Exists (

Select $C_2.Sid$
From Catalog C_2

Where ($C_2.Pid = P.Pid$

AND
 $C_2.Sid = C_1.Sid$)

C_2	
Sid	Pid
S_1	P_1
S_2	P_2
S_1	P_2

- ∴ Not EXISTS will return True
- ∴ P_1 is selected

W.r.t. 1st tuple of parts 'P' inner query result is empty

Select $C_1.Sid$

From Catalog C_1

Where Not Exists

w.r.t 2nd tuple
Catalog (C_1) inner

query result is not empty

o. Not EXISTS will
return false

& Sid of the 2nd tuple
is not selected

C_1	
Sid	Pid
S_1	P_1
S_2	P_2
S_1	P_2

Select
From Parts P

Where Not Exists

P	
Pid	
P_1	
P_2	

- o. Not EXISTS will return false
- o. P_2 is not selected

C_2	
Sid	Pid
S_1	$P_1 X$
S_2	$P_2 V$

Select $C_2.Sid$

From Catalog C_2

Where $(C_2.Pid = P.Pid)$

AND
 $C_2.Sid = C_1.Sid$)

w.r.t 2nd tuple
of parts 'P' inner
query result is
not empty

Select $C_1.Sid$
From Catalog C_1
Where Not Exists (

C_1	
Sid	Pid
S_1	P_1
S_2	P_2
S_1	P_2

Select $P.Pid$
From Parts P

P	
Pid	
P_1	
P_2	

Where Not Exists (Select $C_2.Sid$
From Catalog C_2
Where ($C_2.Pid = P.Pid$
AND
 $C_2.Sid = C_1.Sid$))

C_2	
Sid	Pid
S_1	P_1
S_2	P_2
S_1	P_2

- ∴ Not EXISTS will return false
- ∴ P_1 is not selected

W.r.t. 1st tuple
of parts 'P' inner
query result is
not empty

Select $C_1.Sid$

From Catalog C_1

Where Not Exists

w.r.t 3rd tuple w.r.t Catalog (C_1) inner

query result is empty

o. Not EXISTS will return true

& Sid of the 3rd tuple is selected

C_1	
Sid	Pid
S_1	P_1
S_2	P_2
S_1	P_2

Select
From Parts P

Where Not Exists

(Select $C_2.Sid$
From Catalog C_2
Where $(C_2.Pid = P.Pid$
AND
 $C_2.Sid = C_1.Sid)$)

P	
Pid	
P_1	
P_2	

C_2	
Sid	Pid
S_1	P_1
S_2	P_2
S_1	P_2

- o. Not EXISTS will return false
- o. P_2 is not selected

w.r.t 2nd tuple of parts 'P' inner query result is not empty

Select $C_1.Sid$

From Catalog C_1

Where Not Exists

C_1	
Sid	Pid
S_1	P_1
$\times S_2$	P_2
$\checkmark S_1$	P_2

Select $P.Pid$
From Parts P

P	
	Pid
	P_1
	P_2

Final O/P =

Sid
S_1
S_1

Where Not Exists (Select $C_2.Sid$
From Catalog C_2

C_2	
Sid	Pid
S_1	P_1
S_2	P_2
S_1	P_2

Where ($C_2.Pid = P.Pid$

AND
 $C_2.Sid = C_1.Sid$)

#e.g.

```
SELECT C1.sid  
FROM Catalog C1  
WHERE NOT EXISTS ( SELECT P.Pid
```

Catalog

Sid	Pid
S ₁	P ₁
S ₂	P ₂
S ₁	P ₂

Parts

Pid
P ₁
P ₂

```
FROM Parts P
```

```
WHERE NOT EXISTS (SELECT C2.Sid  
FROM Catalog C2  
WHERE(C2.Pid=P.Pid AND C2.Sid=C1.Sid)))
```

What output is produced by above SQL query:

- A. Sids of suppliers who supplied some parts
- B. Sids of suppliers who supplied only proper subset of parts from all parts
- C. Sids of suppliers who supplied all parts
- D. Sids of suppliers who did not supply any part



Topic : AS clause

- AS clause is used to rename a column or table with an alias. { Note: Using 'AS' we can create an alias }
for almost Everything
- An alias only exists for the duration of the query.

#e.g PYQ

Consider a database that has the relation schema **EMP** (Empld, EmpName, and DeptName).

An instance of the schema EMP and a SQL query on it are given below:

Empld	Emp Name	DeptName
1	XYA	AA
2	XYB	AA
3	XYC	AA
4	XYD	AA
5	XYE	AB
6	XYF	AB
7	XYG	AB
8	XYH	AC
9	XYI	AC
10	XYJ	AC
11	XYK	AD
12	XYL	AD
13	XYM	AE

Main Query

SELECT AVG(EC.Num)

FROM EC

WHERE (DeptName, Num) IN

(SELECT DeptName, COUNT(Empld) AS EC(DeptName, Num)

FROM EMP

GROUP BY DeptName)

aggregate function

The output of executing the SQL query is

Independent Sub-query
:: It will be executed first

#e.g. Q10
GATE-
PYQ

Consider a database that has the relation schema **EMP** (Empld, EmpName, and DeptName).

An instance of the schema EMP and a SQL query on it are given below:

Empld	Emp Name	DeptName
1	XYA	AA
2	XYB	AA
3	XYC	AA
4	XYD	AA
5	XYE	AB
6	XYF	AB
7	XYG	AB
8	XYH	AC
9	XYI	AC
10	XYJ	AC
11	XYK	AD
12	XYL	AD
13	XYM	AE

SELECT AVG(EC.Num)

FROM EC

WHERE (DeptName, Num) IN

SELECT DeptName, COUNT(Empld) AS EC(DeptName, Num)

FROM EMP
GROUP BY DeptName)

The output of executing the SQL query is

name
given to
the table
w.r.t o/p

new
name
w.r.t.
first
attribute

new
name
w.r.t-
2nd
attribute

its a new schema

#e.g. Consider a database that has the relation schema EMP (Empld, EmpName, and DeptName).

An instance of the schema EMP and a SQL query on it are given below:

```
SELECT AVG(EC.Num)
```

```
FROM EC
```

```
WHERE (DeptName, Num) IN
```

```
(SELECT DeptName, COUNT(Empld) AS EC(DeptName, Num)
```

```
FROM EMP
```

```
GROUP BY DeptName)
```

- inner query The output of executing the SQL query is
- it will be executed first.
- And because of the Execution of this query a new Schema "EC(DeptName, Num)" will come into the picture.
- then Main query will execute w.r.t. table EC(Deptname, Num)

#e.g. Consider a database that has the relation schema EMP (Empld, EmpName, and DeptName). An instance of the schema EMP and a SQL query on it are given below:

EC

DeptName	Num
DeptName	Count(EmpID)
A A	4
A B	3
A C	3
A D	2
A E	1

```
SELECT AVG(EC.Num)
FROM EC
WHERE (DeptName, Num) IN
(SELECT DeptName, COUNT(EmpId) AS EC(DeptName, Num)
FROM EMP
GROUP BY DeptName)
```

The output of executing the SQL query is

#e.g.

Consider a database that has the relation schema
EMP (Empld, EmpName, and DeptName).

An instance of the schema EMP and a SQL query on it are given below:

P
W

$$\frac{\text{SUM}(EC.\text{Num})}{\text{Count}(EC.\text{Num})} = \frac{13}{5} = 2.6$$

SELECT AVG(EC.Num)
FROM EC
WHERE (DeptName, Num) IN
(SELECT DeptName, COUNT(Empld) AS EC(DeptName, Num)
FROM EMP
GROUP BY DeptName)

EC

DeptName	Num
AA	4
AB	3
AC	3
AD	2
AE	1

The output of executing the SQL query is _____

Set of tuple produced
by this query is
 $\{(AA, 4), (AB, 3), (AC, 3), (AD, 2), (AE, 1)\}$

#e.g. Consider a database that has the relation schema EMP (Empld, EmpName, and DeptName).

An instance of the schema EMP and a SQL query on it are given below:

```
SELECT AVG(EC.Num)
FROM EC
WHERE (DeptName, Num) IN
    (SELECT DeptName, COUNT(Empld) AS EC(DeptName, Num)
     FROM EMP
     GROUP BY DeptName)
```

The output of executing the SQL query is _____

$$\underline{\text{AM}} = \underline{\underline{2.6}}$$



Topic : WITH clause

WITH is used along with AS

The WITH Clause is mainly used to provide a subquery block a name that can be referenced by the main SQL query or the subqueries that follows.



Syntax of WITH :-

new name
assigned to
O/P of sub-query

New names
assigned to
the Columns

O/P of this subquery will
be recognized by a "new schema"
defined using WITH clause

WITH new-table-name (new
name w.r.t attribute, A₁,
new name w.r.t attribute, A₂, - - -) AS

We may have
multiple
Sub-queries
Using "WITH"
Clause

All of them will execute
in a sequence one after another

Sub-query

O/P =

A ₁	A ₂	-	-

In the
end we
will have
Main
Query

Select - - - - -
- - - - - ← Main query

#e.g. Consider the following database table named water_scheme

GATE
PYQ

water_scheme		
scheme_no	district_name	capacity
1	Ajmer	20
1	Bikaner	10
2	Bikaner	10
3	Bikaner	20
1	Churu	10
2	Churu	20
1	Dungargarh	10

Order of execution

The number of tuples returned by the following SQL query is _____

- ① `with total(name, capacity) as
select district_name, sum(capacity)
from water_schemes
group by district_name`
- ② `with total_avg(capacity) as
select avg(capacity)
from total`
- ③ `select name from total, total_avg
where total.capacity >= total_avg.capacity`

Main Query

#e.g. Consider the following database table named water_scheme

GATE
PYQ

water_scheme		
scheme_no	district_name	capacity
1	Ajmer	20
1	Bikaner	10
2	Bikaner	10
3	Bikaner	20
1	Churu	10
2	Churu	20
1	Dungargarh	10

order of execution

The number of tuples returned by the following SQL query is _____

④ with total(name, capacity) as
 Subquery { ③ select district_name, sum(capacity)
 ① from water_schemes
 ② group by district_name
 ⑦ with total_avg(capacity) as
 Subquery { ⑥ select avg(capacity)
 ⑤ from total
 ⑩ select name from total, total_avg
 ⑧ where total.capacity >= total_avg.capacity
 ⑨

#e.g. Consider the following database table named water_scheme

district_name	Sum(Capacity)
Ajmer	20
Bikaner	40
Churu	30
Dungarpur	10

The number of tuples returned by the following SQL query is _____

```
with total(name, capacity) as
    select district_name, sum(capacity)
        from water_schemes
            group by district_name
with total_avg(capacity) as
    select avg(capacity)
        from total
select name from total, total_avg
    where total.capacity >= total_avg.capacity
```

#e.g. Consider the following database table named water_scheme

total

Name	Capacity
district_name	Sum(Capacity)
Ajmer	20
Bikaner	40
Churu	30
Dungarpur	10

total_avg

Capacity
Avg(Capacity)
25

The number of tuples returned by the following SQL query is _____

```
with total(name, capacity) as
  select district_name, sum(capacity)
  from water_schemes
  group by district_name
```

```
with total_avg(capacity) as
  select avg(capacity)
  from total
```

```
select name from total, total_avg
  where total.capacity >= total_avg.capacity
```

#e.g. Consider the following database table named water_scheme

total	
Name	Capacity
Ajmer	20
Bikaner	40
Churu	30
Dungarpur	10

total_avg	
Capacity	
25	

The number of tuples returned by the following SQL query is _____

```
with total(name, capacity) as
  select district_name, sum(capacity)
  from water_schemes
  group by district_name
with total_avg(capacity) as
  select avg(capacity)
  from total
select name from total, total_avg
  where total.capacity >= total_avg.capacity
```

O/P =	
name	
Bikaner	
Churu	

⇒ # tuples in O/P = 2



2 mins Summary



- Topic** Execution of correlated sub-query
- Topic** AS clause
- Topic** WITH clause

THANK - YOU