

CS & IT ENGINEERING



Operating System

File System

(One Shot)



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Recap of Previous Lecture



Topic

Inverted Paging

Topic

Hashed Page table

Topics to be Covered



Topic

File System

Topic

Blocks and File Allocation Methods

Topic

Disk Scheduling

Topic

SSTF, Scan, Look



Topic : File

A file is a named collection of related information that is recorded on secondary storage.





Topic : File Attributes

1. Name
2. Extension
3. Size
4. Date
5. Author
6. Created, Modified, Accessed
7. Attributes: Read-only, hidden
8. Default Program
9. Security Details



Topic : File Directory

→ folder

Collection of files



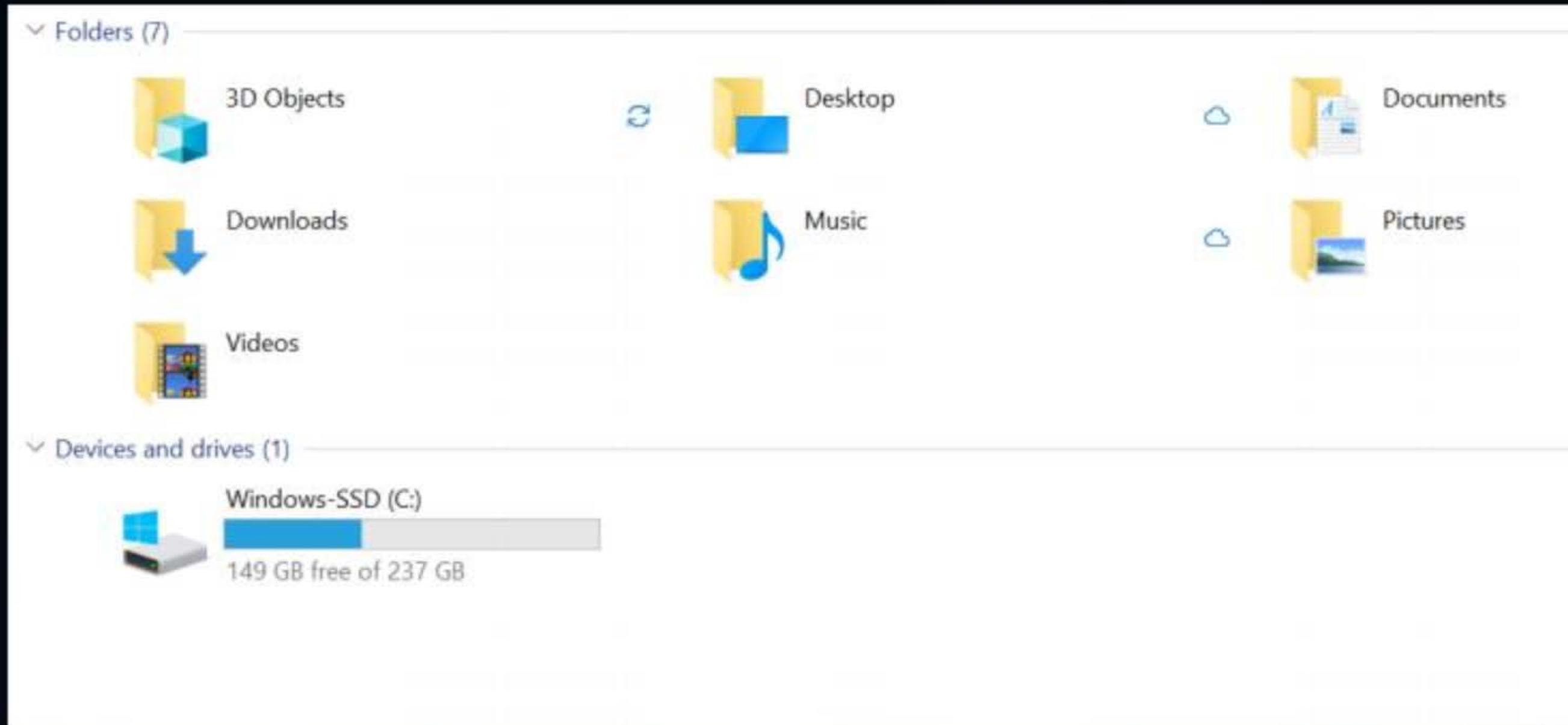
Topic : File System

Module of OS which manages, controls and organizes files and related structures





Topic : File System



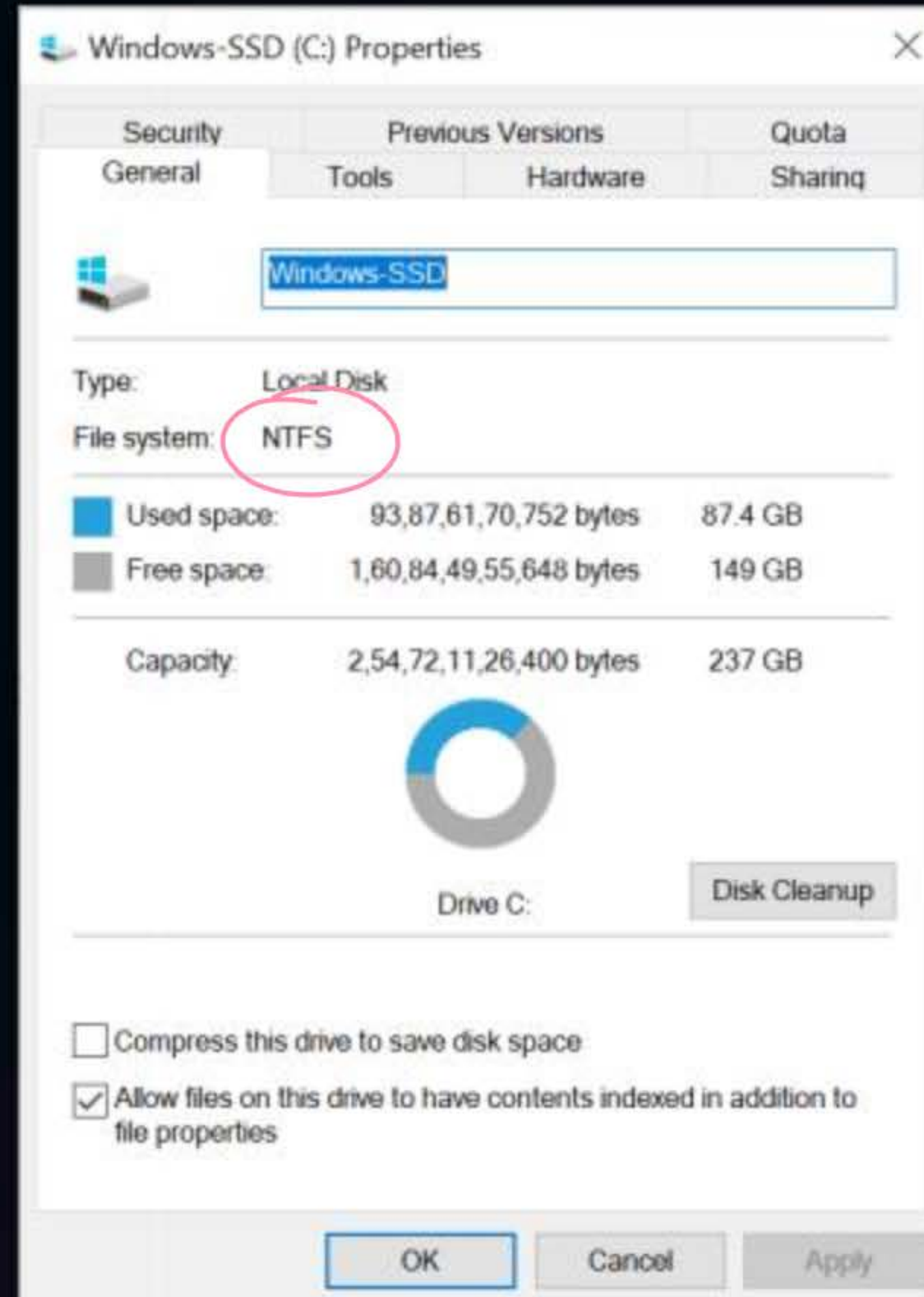


Topic : Types of File Systems

1. FAT32
2. NTFS → windows
3. HFS+ → iOS or macOS
4. Ext2 / Ext3 / Ext4 → linux / unix
5. Swap ✓



Topic : File System





Topic : Disk Formatting

Low level formatting
(physical)

creating tracks and sectors
on disk.

Done by manufacturer

High level formatting
(logical)

⇒ creating drives (partitions)
on disk like C:\ E:\ F:\

→ creating blocks
⇒ 2 types of partitions

Primary



OS + user
files

extended

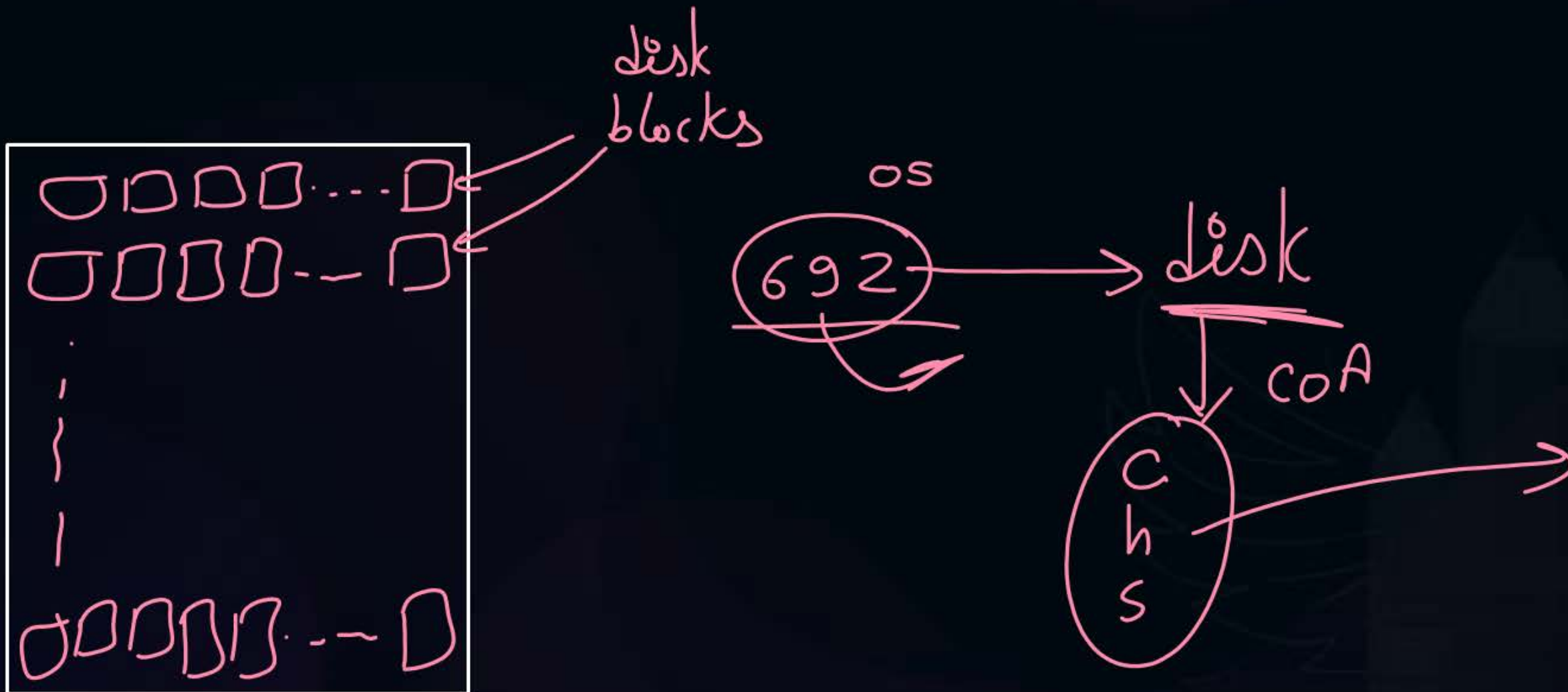


user files



Topic : Disk Blocks

smallest logical unit of disk (1 sector \Rightarrow 1 block)





Topic : Disk Blocks



Number of disk blocks = 2^{16}

Size of each block = 1KB

Total Size of disk?

\Rightarrow Disk block address (D.B.A.) = 16 bits

$$2^{16} * 1KB = 2^{26} B = 64MB$$





Topic : Disk Blocks



Disk block address = 24-bits \Rightarrow no. of blocks = 2^{24}

Size of each block = 2KB

Total Size of disk? = $2^{24} * 2KB = 2^{35} B = 32 GB$





Topic : Disk Blocks



Total disk size = 256GB

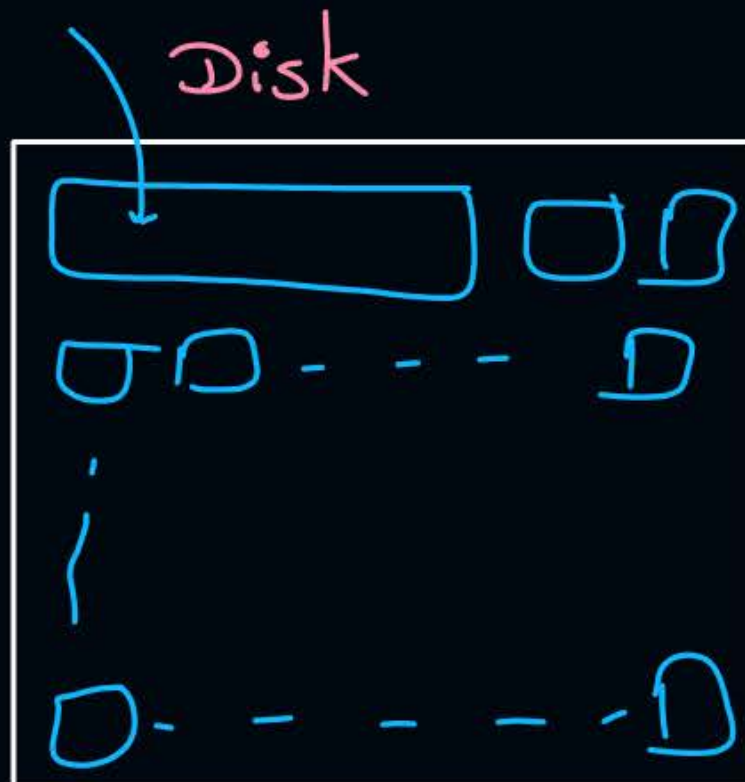
Block Size = 2KB

Disk block address?

$$\text{no. of blocks} = \frac{256 \text{ GB}}{2 \text{ KB}} = \frac{2^{38}}{2^{11}} = 2^{27}$$

$$\text{D.B.A.} = \underline{\underline{27 \text{ bits}}}$$

some blocks are used to store file directory infoⁿ



which file stored where and how

ex:-

A disk has 256 blocks each of size 2KB.
starting 10 blocks are used for storing
directory entries.

max size of disk remaining to store files ?

solⁿ blocks remaining for file = $256 - 10$
 $= 246$

Size for file = $246 * 2KB = 492 KB$

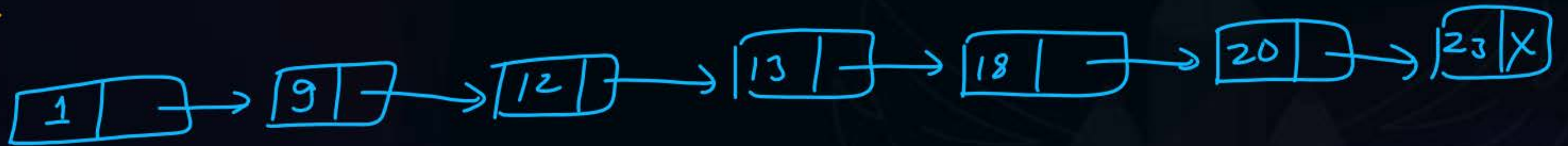


Topic : Free Space Management

1. Free List (linked list of block numbers of all free blocks)
2. Bitmap Method (one bit per block)

ex:- blocks 1, 9, 12, 13, 18, 20, 23 free

free list:-



Bitmap:- $\begin{cases} 0 \Rightarrow \text{occupied} \\ 1 \Rightarrow \text{free} \end{cases}$

Block add starts from 0 \Rightarrow 010000000100110000101001



Topic : Free Space Management



1. No searching in free list, but in bitmap we search for first ~~free~~ free block
2. Free list is faster in allocating a free block
3. Free list size is variable, where as bitmap size is constant





Topic : Question

A particular disk unit uses a bit string to record the occupancy or vacancy of its disk blocks with '0' denoting vacant block and '1' denoting occupied block. A 32-bit part of this string has Hexadecimal value of A3D4B055. The percentage of occupied blocks on the disk for this part is ?

$$\begin{array}{cccccccc} 1010 & 0011 & 1101 & 0100 & 1011 & 0000 & 0101 & 0101 \\ \hline 2 & 2 & 3 & 1 & 3 & 0 & 2 & 2 \end{array}$$

$$\frac{15}{32} * 100\% = 46.875\%$$

Ques) A disk has block size = 2KB

$$\text{no of blocks} = 2^{24}$$

$$\text{no. of blocks needed to store bitmap for free blocks} = \underline{1024} ?$$

Solⁿ

$$\begin{aligned} \text{no. of bits in bitmap} &= \text{no. of blocks} \\ &= 2^{24} \text{ bits} \\ &= 16 \text{ M bits} \end{aligned}$$

$$\text{no. of blocks} = \frac{16 \text{ M bits}}{2 \text{ KB}}$$

$$= \frac{2 \text{ MB}}{2 \text{ KB}}$$

$$= 1 \text{ K}$$

$$= 1024$$



Topic : File Allocation Method ✓



↓
describes which disk blocks to be allocated to a file.





Topic : File Allocation Method

1. Contiguous Allocation
2. Linked Allocation
3. Indexed Allocation



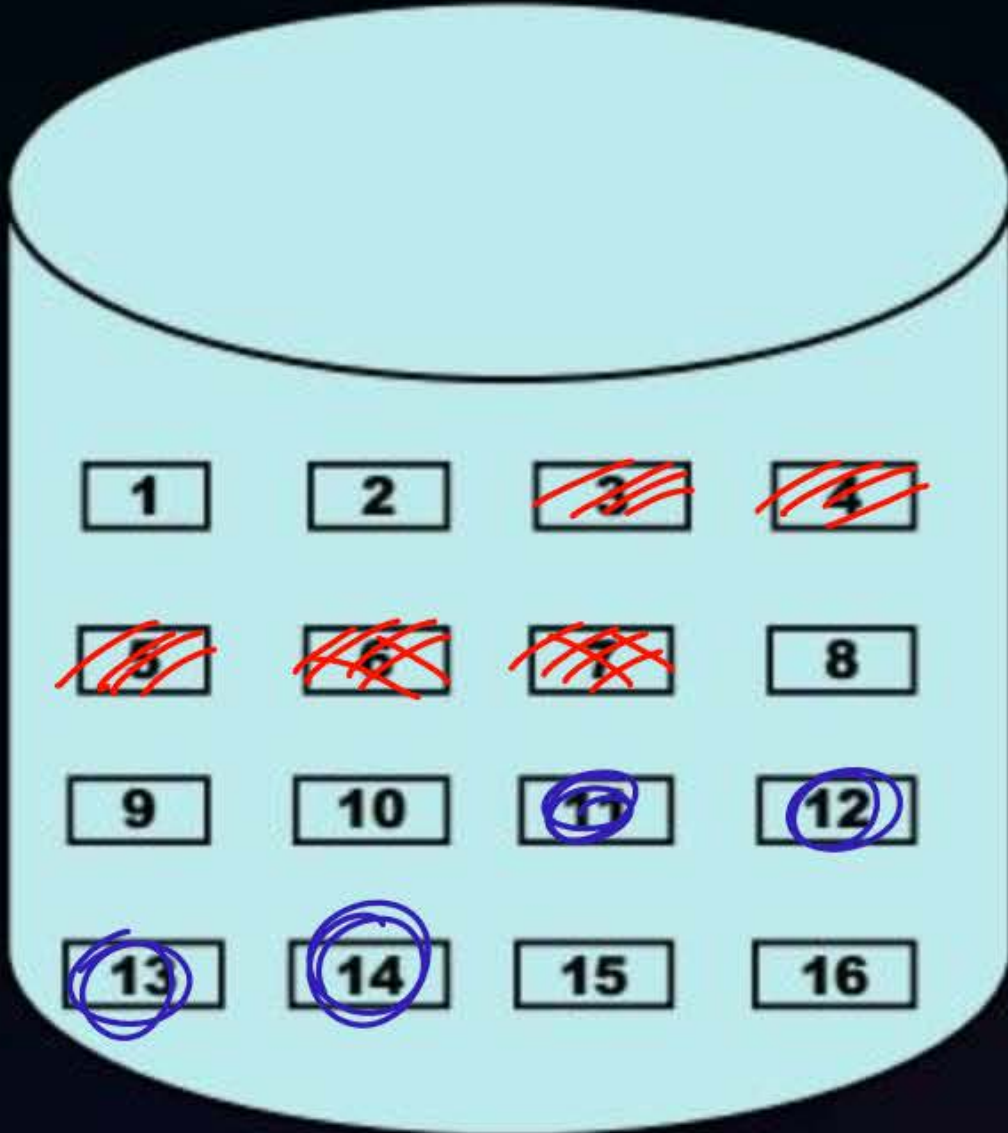
Topic : Contiguous Allocation



consecutive blocks are used to store file

file allocation

file name	starting block no.	no. of blocks
os.pdf	3	5
coA.pptx	11	4





Topic : Contiguous Allocation

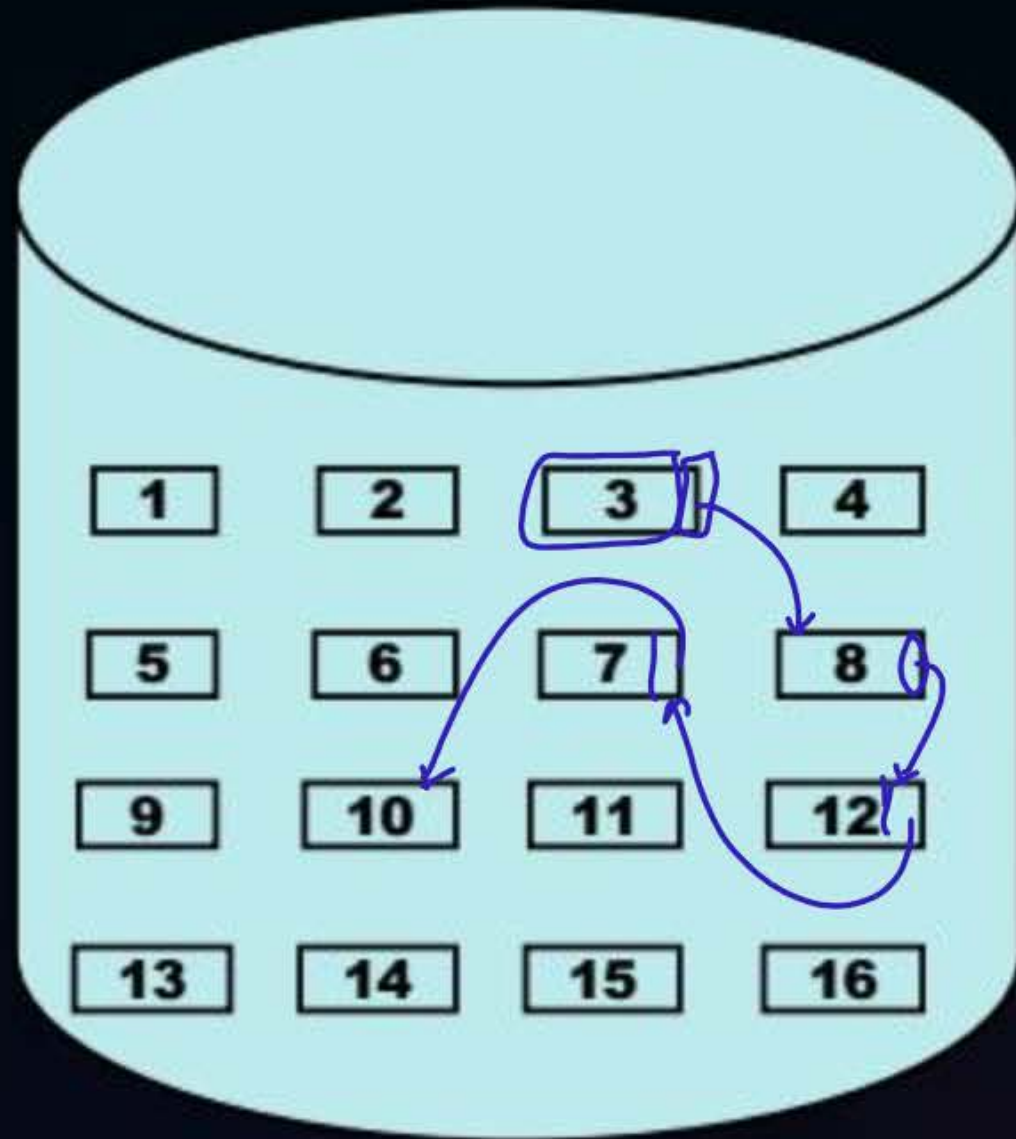
Performance:

1. Fragmentation: Internal, External
2. Increase in File size: Inflexible
3. Type of access: Sequential, Random/direct

Insertion of a new block in file at	no. of block accesses
Beginning	many block accesses needed
Middle	———— —————
End	easy and lesser block accesses



Topic : Linked Allocation



file name	starting block no.	last block no	no. of blocks 5
coA. pptx	3	10	



Topic : Linked Allocation

Performance:

1. Fragmentation: Internal
2. Increase in File size: Flexible
3. Type of access: Sequential

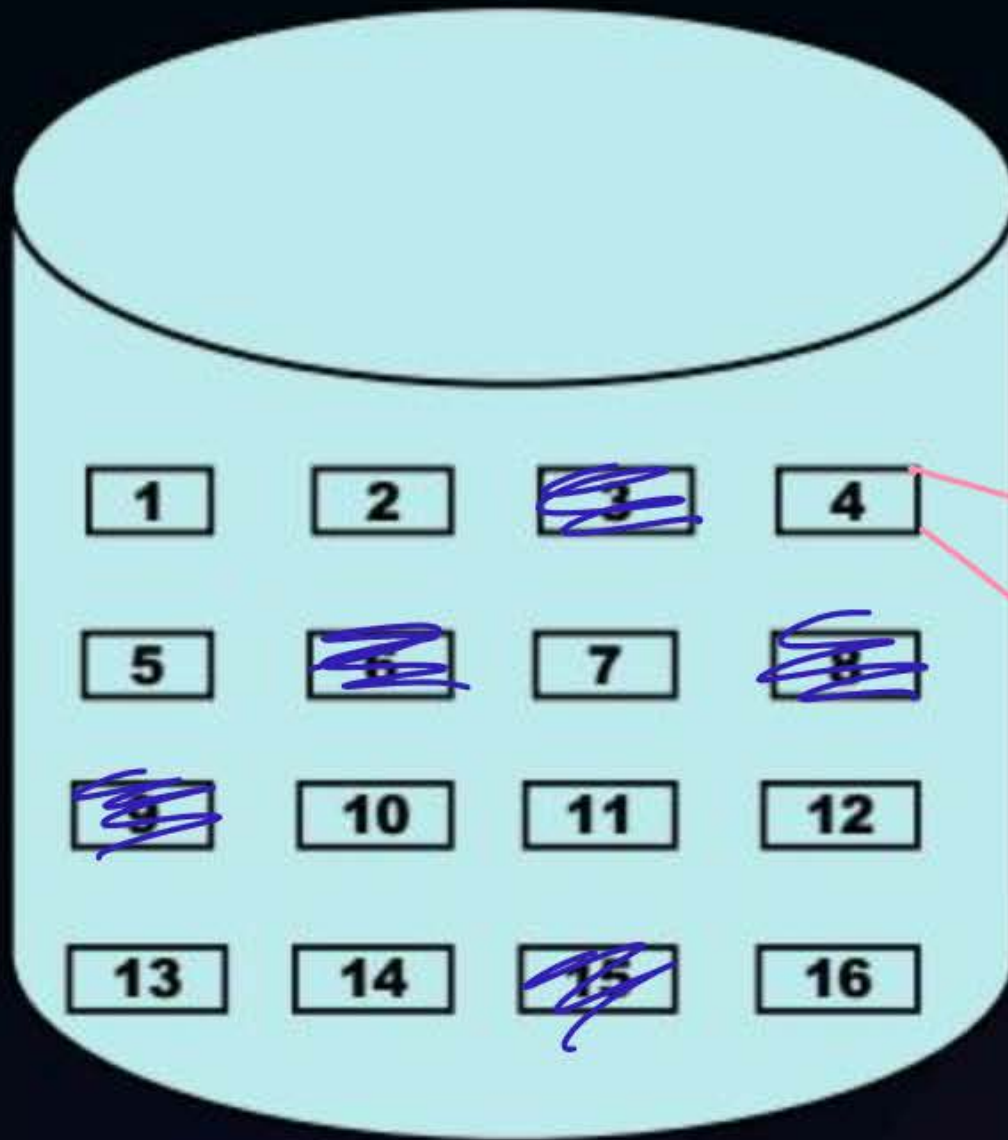
Insertion of a new block in file at	no. of blocks accessed
Beginning	very less
Middle	more (as traversal needed)
End	more (— —)



Topic : Indexed Allocation

Allocation table

file name	Index block no.
COA.pptx	4



Index block

6
9
3
8



Topic : Indexed Allocation

Performance:

1. Fragmentation: Internal
2. Increase in File size: Flexible
3. Type of access: Sequential, Random/direct

If file is very big
then
use multilevel
indexing

Insertion of a new block in file at	no. of block accesses
Beginning	very less
Middle	- -
End	- -



Topic : Question

H. ω.

[GATE-2022]



#Q. Consider two files systems A and B , that use contiguous allocation and linked allocation, respectively. A file of size 100 blocks is already stored in A and also in B. Now, consider inserting a new block in the middle of the file (between 50th and 51st block), whose data is already available in the memory. Assume that there are enough free blocks at the end of the file and that the file control blocks are already in memory. Let the number of disk accesses required to insert a block in the middle of the file in A and B are n_A and n_B respectively, then the value of $n_A + n_B$ is_____?



Topic : Question

- Disk block address = 16 bits = $2B$
- Disk block size = 1KB
- Index block = 1KB
- Maximum file size? (single level indexing)

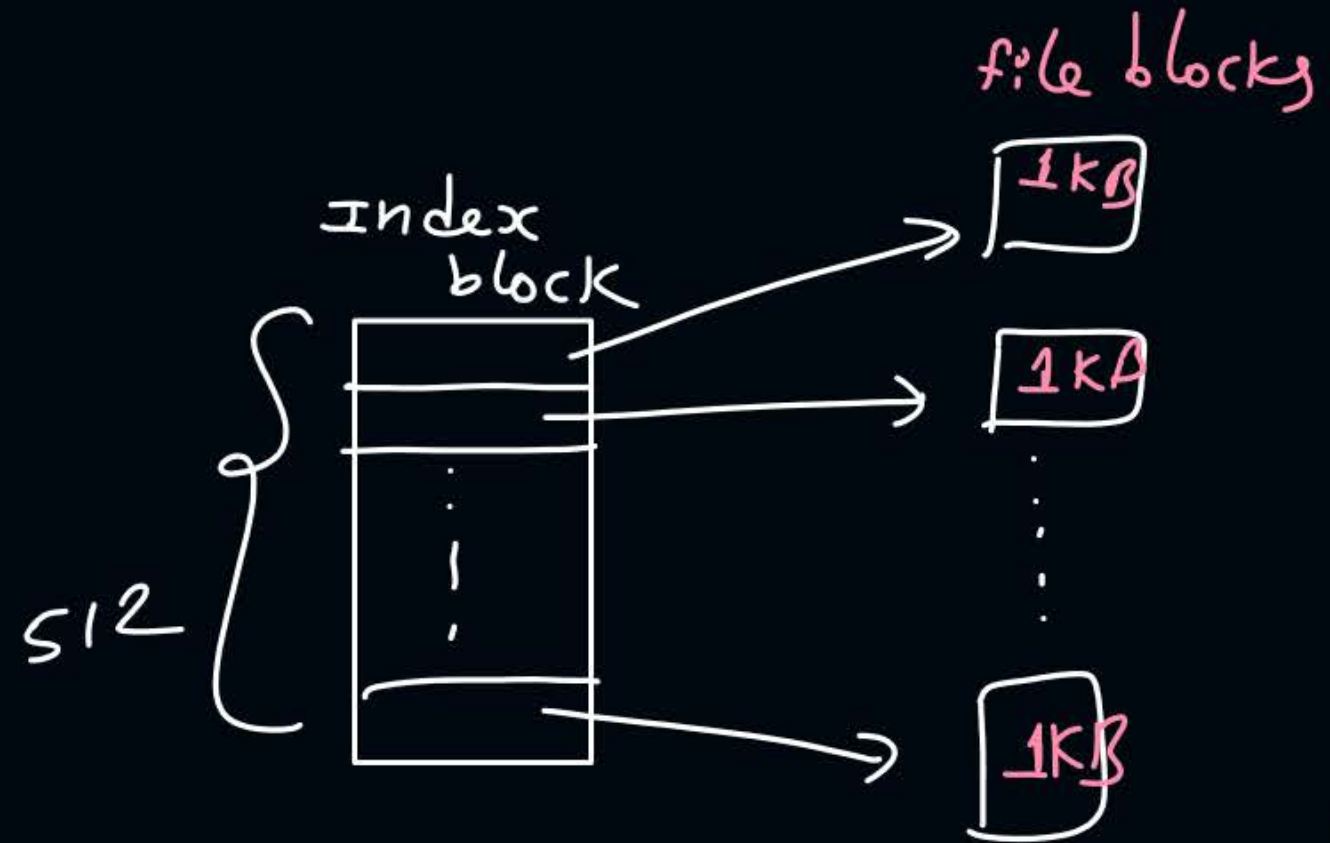
$$\text{no. of disk block numbers per block} = \frac{1KB}{2B}$$

$$= 512$$

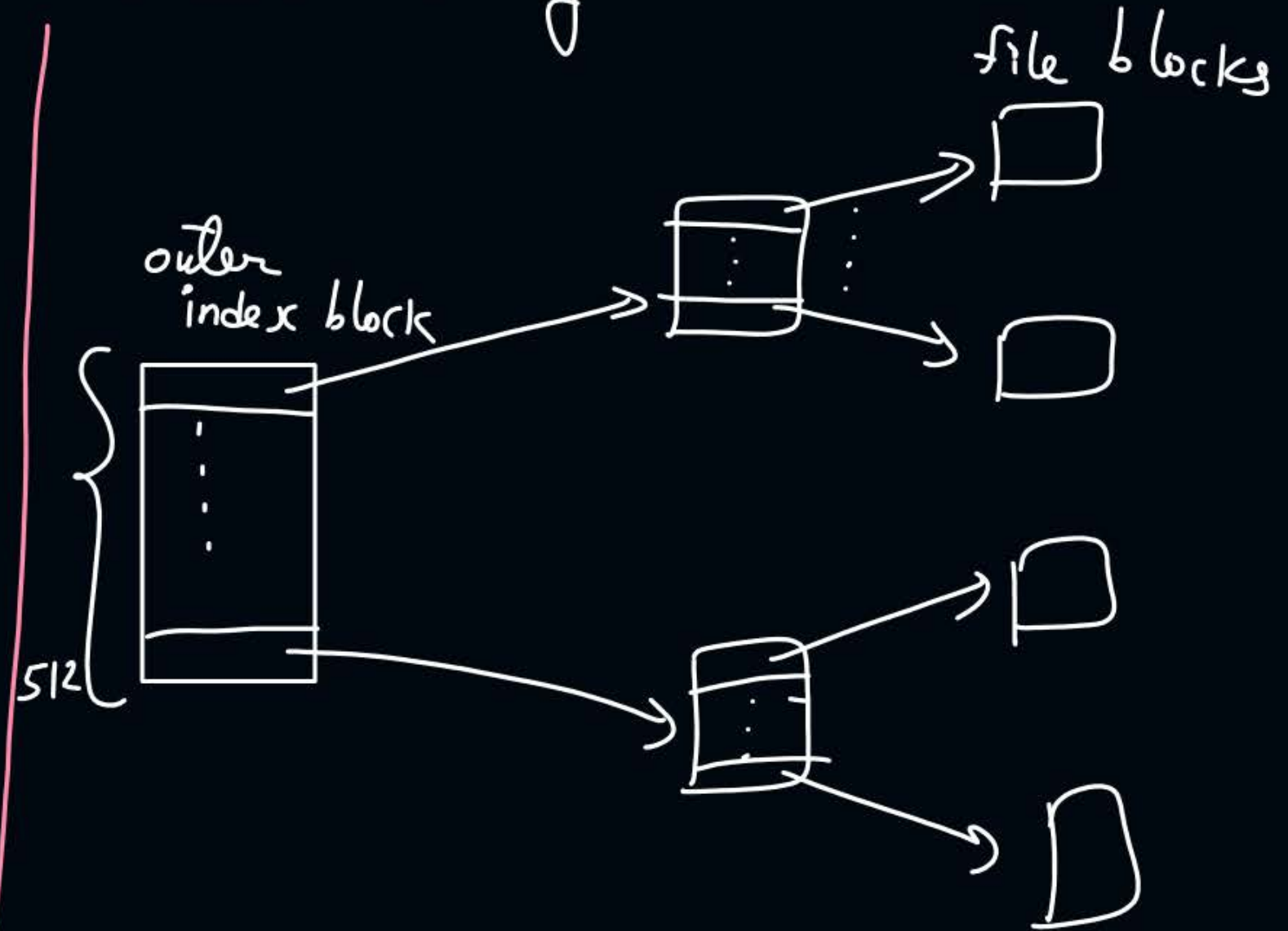
$$\text{max file size} = 512 * 1KB$$

$$= 512KB$$

1-level indexing



2-level indexing





Topic : Question

- Disk block address = 16 bits = 2^{16}
- Disk block size = 1KB
- Index block = 1KB
- Maximum file size? (double level indexing)

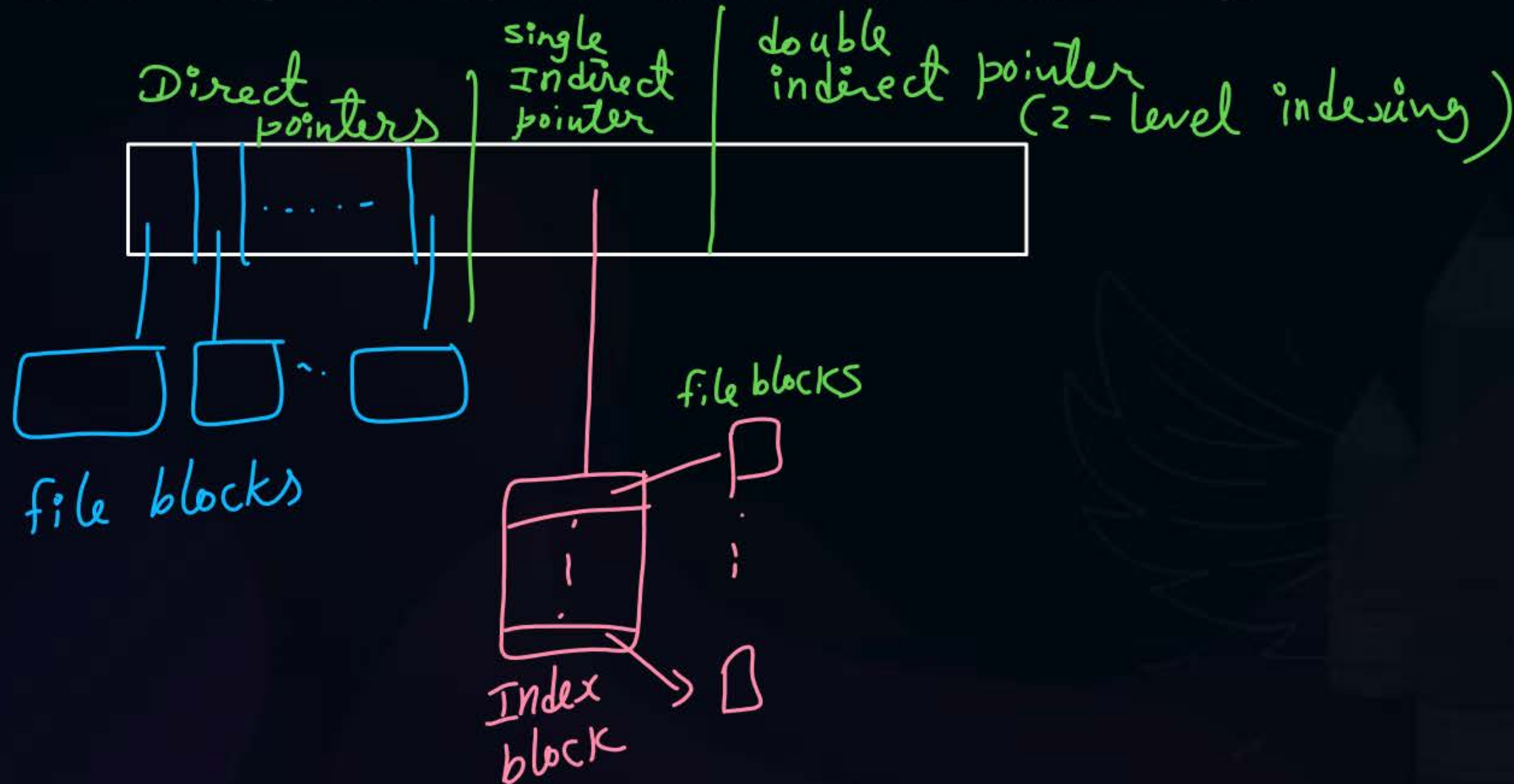
$$\text{no. of disk block numbers per block} = \frac{1\text{KB}}{2^{16}}$$
$$= 512 = 2^9$$

$$\text{max file size} = 2^9 * 2^9 * 1\text{KB} = 256\text{MB}$$



Topic : Unix I-node Structure

The inode (index node) is a data structure in a Unix-style file system that describes a file-system object such as a file or a directory.





Topic : Unix I-node Structure

The inode (index node) is a data structure in a Unix-style file system that describes a file-system object such as a file or a directory.

- Each inode stores the attributes and disk block locations of the object's data.
- The number of Inode limits the total number of files/directories that can be stored in the file system.



Topic : Question

[GATE-2019]



#Q. The index node (inode) of a Unix-like file system has 12 direct, one single-indirect and one double-indirect pointer. The disk block size is 4 kB, and the disk block addresses 32-bits long. The maximum possible file size is (rounded off to 1 decimal place) 4.0 GB?

4B

$$\text{no. of indexes per block} = \frac{4 \text{ KB}}{4 \text{ B}} = 1 \text{ K}$$

$$\begin{aligned} \text{max file size} &= (12 * 4 \text{ KB}) + (1 \text{ K} * 4 \text{ KB}) + (1 \text{ K} * 1 \text{ K} * 4 \text{ KB}) \\ &= 48 \text{ KB} + 4 \text{ MB} + 4 \text{ GB} \\ &\approx 4.004048 \text{ GB} = 4.0 \text{ GB} \end{aligned}$$



Topic : Question

H.W.

[GATE-2005]



#Q. In a computer system, four files of size 11050 bytes, 4990 bytes, 5170 bytes and 12640 bytes need to be stored. For storing these files on disk, we can use either 100 byte disk blocks or 200 byte disk blocks (but can't mix block sizes). For each block used to store a file, 4 bytes of bookkeeping information also needs to be stored on the disk. Thus, the total space used to store a file is the sum of the space taken to store the file and the space taken to store the book keeping information for the blocks allocated for storing the file. A disk block can store either bookkeeping information for a file or data from a file, but not both.

What is the total space required for storing the files using 100 byte disk blocks and 200 byte disk blocks respectively?

A

35400 and 35800 bytes

B

35800 and 35400 bytes

C

35600 and 35400 bytes

D

35400 and 35600 bytes



Topic : Solution

File Size	No. of blocks to store file	No. of blocks to store bookkeeping	Total blocks	Total Size
11050 Bytes				
4990 Bytes				
5170 Bytes				
12640 Bytes				



Topic : Solution

File Size	No. of blocks to store file	No. of blocks to store bookkeeping	Total blocks	Total Size
11050 Bytes				
4990 Bytes				
5170 Bytes				
12640 Bytes				



2 mins Summary

Topic

File System

Topic

Blocks and File Allocation Methods

Topic

Disk Scheduling

Topic

SSTF, Scan, Look



Happy Learning

THANK - YOU

