

Class Notes

CS & IT ENGINEERING



Theory of Computation

Turing Machine
Undecidability

DPP_01

By- Rohit Chauhan sir



[MCQ]



#Q. Consider the following two regular expressions R_1 and R_2 over $\Sigma = \{a, b\}$.

✓ $R_1 = a^* (ba^*)^* \Rightarrow (a+b)^*$

✓ $R_2 = (a^* + b^*)^* = (a+b)^*$

Which of the following is true?

$$\frac{a^* (ba^*)^*}{(ba^*)^2} = \{ \epsilon, a^*, b^*, a^*b^*, b^*a^*, \dots \}$$

$$\begin{array}{l} ba^*ba^* \\ \hline bba^* \end{array}$$

A $R_1 \subset R_2$

B $R_2 \subset R_1$

C $R_1 = R_2$

D $R_1 \cap R_2 = a^*b^*$



#Q. Consider the following statements:

- ✓ I. A regular grammar can generate some CFL language. ✓
- ✓ II. There are some non-regular language for which PDA not exist. ✓
- ✓ III. Context sensitive grammar exists for every CFL. ✓
- ✓ IV. If L is a regular and M is not a regular language then L.M is necessarily not regular. ✓

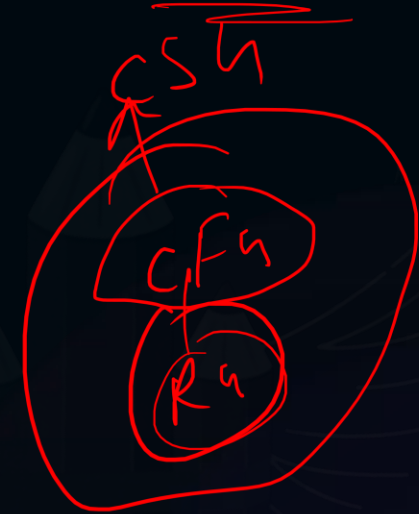
Number of correct statement is/are _____?

3

$\{a^n b^n / n \geq 0\}$ NR \rightarrow PDA
 $\{a^n b^m c^n / n \geq 0\}$ CSL \rightarrow Non-Reg

CSL \rightarrow LBA

$\phi \cdot \{a^n b^n / n \geq 0\}$
 $\Rightarrow \phi$



#Q. Which of the following languages are not decidable?

A

$L = \{ \langle G \rangle \mid G \text{ is CFG and } R \text{ is a regular set such that } L(G) = R \}$

UD

CFL

B

$L = \{ \langle G \rangle \mid G \text{ is CFG such that } L(G) = \phi \}$

Decidable

mem ✓

C

$L = \{ \langle G_1, G_2 \rangle \mid G_i \text{ is CFG such that } L(G_1) \cap L(G_2) = \Sigma^* \}$

completeness ✓
Finiteness ✓
Emp ✓

D

$L = \{ \langle G_1, G_2 \rangle \mid G_i \text{ is CFG such that } L(G_1) = L(G_2) \}$

UD

UD

#Q. Consider the following problems:

P_1 : Does a given program ever produce an output? \rightarrow $\forall \emptyset$

P_2 : Given DFAs D_1 and D_2 is $L(D_1) \cap L(D_2) = \emptyset$? Decidable

P_2 : For Context sensitive grammar G and string w , is $w \in L(G)$? membership

How many languages is/are Undecidable? 1.

Non-Trivial

1

Decidable

[MCQ]



#Q. If $L_1 \cap L_2$ is regular and L_1 is non-regular then, L_2 must be

$$\left. \begin{array}{l} L_1 \cap L_2 = \text{Reg} \\ \text{NR} \cap L_2 = \text{Reg} \end{array} \right\} \begin{array}{l} \{a^n b^n\} \cap \phi = \phi \\ \phi = \text{Reg} \\ a b = \text{Reg} \end{array} \quad \textcircled{D}$$

A Regular but not finite ~~X~~

B Non-regular ~~X~~

C Finite ~~X~~

~~**D** None of these~~

$$\underbrace{\{a^n b^n \mid n \geq 0\}}_{L_1} \cap \underbrace{\{a^n b^n \mid n \geq 0\}}_{\phi}$$

$$P_1 \leq P_2$$

#Q. If P_1 is reducible to P_2 then, which of the following is/are correct?

A

If P_1 is decidable then, P_2 is undecidable. *we can't say X*

B

If P_2 is undecidable then, P_1 is decidable. *X*

C

If P_2 is decidable then, P_1 must be decidable. *✓*

D

If P_1 is undecidable then, P_2 must be undecidable. *✓*

[MSQ]



#Q. Let,

✓ $L_1 = \text{CFL}$

✓ $L_2 = \text{DCFL}$

✓ $L_3 = \text{Regular}$

Then, which of the following is/are correct?

B, D

$$\underline{L_1 - L_2} = L_1 \cap \overline{L_2}$$

$$a^n b^n \cdot \phi = R$$

$$\underline{a^n b^n \cdot c} = \text{DCFL}$$

A

$L_3 - L_1$ is CSL

$$\Rightarrow \text{Reg} \cap \overline{\text{CFL}} \Rightarrow \text{Reg} \cap \text{CSL} \Rightarrow \underline{\text{CSL}}$$

B

$L_1 \cup (L_2 \cap \overline{L_3})$ is CFL

$$\Rightarrow \text{CFL} \cup (\text{DCFL} \cap \overline{\text{Reg}}) = \text{CFL} \cup \text{DCFL} = \text{CFL}$$

C

$L_2 \cdot L_3$ is Regular

$$= \underline{\text{DCFL}} \cdot \text{Reg} = \text{X}$$

D

$L_2 \cup L_3$ is DCFL

$$\text{DCFL} \cup \text{Reg} = \text{DCFL}$$

#Q. Consider the following statements:

S1: Complement of finite language may be finite. ~~X~~

S2: Kleene star of finite language may be finite. ✓

S3: Subset of finite language is always finite. ✓

S4: Complement of infinite language may be finite. ✓

Number of INCORRECT statements is/are ____.

1

$$\begin{aligned}\overline{\Phi} &= (a+b)^* \\ \epsilon &= \overline{(a+b)^* - \epsilon} \\ S2 \Rightarrow \Phi^* &= \epsilon \\ \epsilon^* &= \epsilon\end{aligned}$$

{ ab, ba, aa, bbb }
✓✓

B, D

#Q. Which of the following problems are undecidable?

A

~~Membership problem in context-free languages.~~

B

~~Whether a given context-free language is regular.~~

C

~~Whether a finite state automation halts on all inputs.~~

D

~~Membership problem for type-0 languages.~~

CFL

$\Rightarrow M$

$\Rightarrow F$

$\Rightarrow D$

$\Rightarrow E$

RE

R_n

CSL

CFL

R_0

FSM

Type-0

Type-1

Type-2

Type-3

RE

[MCQ]



#Q. Which one of the following is UNDECIDABLE?

B

- ☒ **A** *Decidable* Given a Turing machine M, a string s and an integer k, M accepts s within k steps
- ☒ **B** Equivalence of two given Turing machines *UD*
- ☒ **C** Language accepted by a given finite state machine is not empty *Decidable*
- ☒ **D** Language generated by a context free grammar is non empty *Decidable*

[MCQ]



#Q. Consider the following decision problems:

(P_1) Does a given finite state machine accept a given string? ✓

(P_2) Does a given context free grammar generate an infinite number of strings

Which of the following statements is true?

A

Both (P_1) and (P_2) are decidable ✓

B

Neither (P_1) nor (P_2) are decidable

C

Only (P_1) is decidable

D

Only (P_2) is decidable

[MCQ]



#Q. Which of the following problems is undecidable?

A Membership problem for CFGs.

Dec

B Ambiguity problem for CFGs.

UD

C Finiteness problem for FSAs.

Dec

D Equivalence problem for FSAs.

Dec

<u>CFL</u>	<u>RM</u>	<u>CSL</u>
F	✓	X
M	✓	✓
E	✓	X

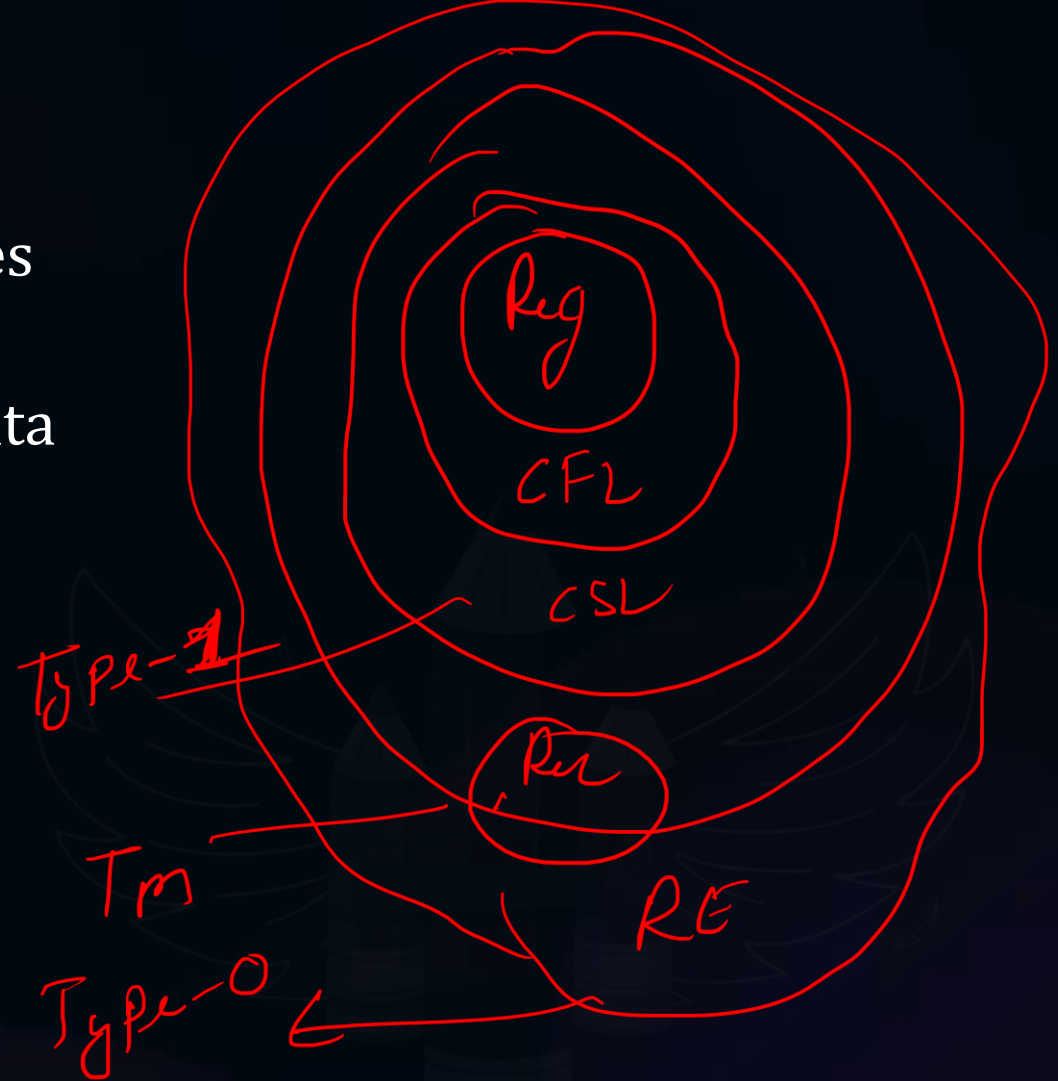
#Q. Recursive languages are:

A a proper superset of context free languages

B ~~always~~ recognized by a pushdown automata

C also called type-0 languages

D recognizable by Turing machines



Classes over

Telegram

SWT \Rightarrow 50-60 problems
 3DPP \Rightarrow 40-50 problems
 CN \Rightarrow 200 problems
 \Rightarrow 100 problem

THANK - YOU

4-6 \rightarrow Reg \rightarrow 50% 60% { Closure properties }

2-4 \rightarrow CFG \rightarrow 25-30%

TM \Rightarrow 1 2 marks

300
250 PYQ

550
DFA \rightarrow RC, RS