

Class Notes

CS & IT ENGINEERING



Theory of Computation

Turing Machine
Undecidability

DPP_01



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#Q. Consider the following two regular expressions R_1 and R_2 over $\Sigma = \{a,b\}$.

$$\checkmark R_1 = a^* (ba^*)^* \Rightarrow (a+b)^*$$

$$\checkmark R_2 = (a^* + b^*)^* = (a+b)^*$$

Which of the following is true?

$$a^* (ba^*)^* = \{(a^*, b^*, ba^*, a^*b^*, b^*a^*, \dots)\}$$

A $R_1 \subset R_2$

B $R_2 \subset R_1$

C $R_1 = R_2$

D $R_1 \cap R_2 = a^*b^*$





[NAT]



#Q. Consider the following statements:

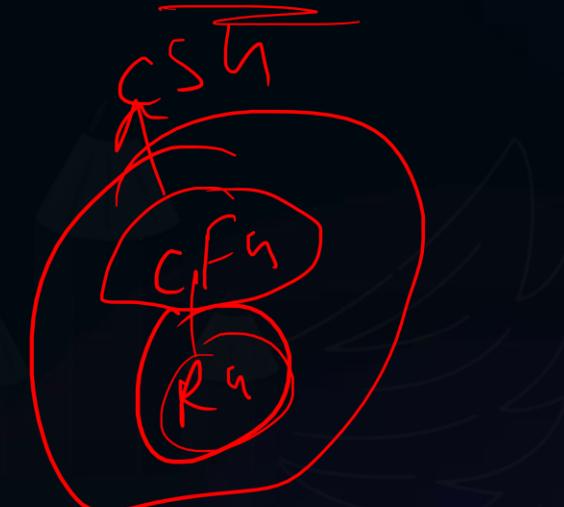
- I. A regular grammar can generate some CFL language. ✓
- II. There are some non-regular language for which PDA not exist. ✓
- III. Context sensitive grammar exists for every CFL. ✓
- IV. If L is a regular and M is not a regular language then L.M is necessarily not regular. ✗

Number of correct statement is/are ____?

$$\begin{array}{c} \{a^nb^n \mid n \geq 0\} \\ \text{NR} \rightarrow \text{PDA} \\ \{a^n b^n \mid n \geq 0\} \\ \text{Non Reg} \end{array}$$

CSL → LBA

∅. {aⁿbⁿ | n ≥ 0} ⇒ ∅



#Q. Which of the following languages are not decidable?

A

$L = \{<G> \mid G \text{ is CFG and } R \text{ is a regular set such that } L(G) = R\}$

NDuc

CFL

B

$L = \{<G> \mid G \text{ is CFG such that } L(G) = \emptyset\}$

Decidable

mem ✓

C

$L = \{<G_1, G_2> \mid G_i \text{ is CFG such that } L(G_1) \cap L(G_2) = \Sigma^*\}$

Finiteness

Emp ✓

Completeness

D

$L = \{<G_1, G_2> \mid G_i \text{ is CFG such that } L(G_1) = L(G_2)\}$

FUD

VO

#Q. Consider the following problems:

P₁: Does a given program ever produce an output? $\rightarrow \text{NP}$

P₂: Given DFAs D₁ and D₂ is L(D₁) \cap L(D₂) = \emptyset ? Decidable

P₃: For Context sensitive grammar G and string w, Is w \in L(G)? Membership

How many languages is/are Undecidable? 1.

N.n Trivial

1

Decidable

[MCQ]



#Q. If $L_1 \cap L_2$ is regular and L_1 is non-regular then, $\underline{L_2}$ must be

$$L_1 \cap L_2 = \text{Reg}$$

$$\text{NR} \cap L_2 = \text{Reg}$$

$$\{a^n b^n\}$$

$$\cap \emptyset =$$

Rg

$$\xleftarrow[a]{ab} = \text{Reg}$$

$$\textcircled{0}$$

A

Regular but not finite



B

Non-regular



C

Finite



None of these

$$\{a^n b^n \mid n \geq 0\}$$

L_1

$$\cap \{a^n b^n \mid n \geq 0\}$$

$$\emptyset$$

$$P_1 \leq P_2$$

#Q. If P_1 is reducible to P_2 then, which of the following is/are correct?

A

If P_1 is decidable then, P_2 is undecidable. *We can't say* \times

B

If P_2 is undecidable then, P_1 is decidable. \times

C

If P_2 is decidable then, P_1 must be decidable.

D

If P_1 is undecidable then, P_2 must be undecidable. \checkmark

[MSQ]

#Q. Let,

$$\checkmark L_1 = \text{CFL}$$

$$\checkmark L_2 = \text{DCFL}$$

$$\checkmark L_3 = \text{Regular}$$

(B, D)

Then, which of the following is/are correct?

A

$$L_3 - L_1 \text{ is CSL} \Rightarrow \text{Reg} \cap \overline{\text{CFL}} \Rightarrow \text{Reg} \cap \text{CSL} \Rightarrow \underline{\text{CSL}}$$

B

$$L_1 \cup (L_2 \cap \overline{L}_3) \text{ is CFL} \Rightarrow \text{CFL} \cup (\text{DCFL} \cap \overline{\text{Reg}}) = \text{CFL} \cup \text{DCFL} = \text{CFL}$$

C

$$L_2 \cdot \overline{L}_3 \text{ is Regular} = \underline{\text{DCFL}} \cdot \text{Reg} = \times$$

D

$$L_2 \cup L_3 \text{ is DCFL} = \text{DCFL} \cup \text{Reg} = \text{DCFL}$$

$$L_1 - L_2 = L_1 \cap \overline{L_2}$$

$$a^{n_1 n} \cdot \phi = R$$

$$\underline{a^{n_1 n} \cdot C} = \text{DCFL}$$

#Q. Consider the following statements:

- { S1: Complement of finite language may be finite. ✗
S2: Kleene star of finite language may be finite. ✓
S3: Subset of finite language is always finite. ✓
S4: Complement of infinite language may be finite. ↗

Number of INCORRECT statements is/are ____.

1

$$\begin{aligned}\overline{\emptyset} &= (\bar{a} + \bar{b})^* \\ \bar{L} &= \overline{(a+b)^* - L} \\ \Rightarrow \emptyset^* &= L \\ S_2 \Rightarrow \emptyset^* &= L\end{aligned}$$

$$\left\{ \overline{ab}, \overline{ba}, \overline{aa}, \overline{bb} \right\}$$

[MSQ]

#Q. Which of the following problems are undecidable?

B, D

A

Membership problem in context-free languages. D

B

Whether a given context-free language is regular. DCFL

C

Whether a finite state automation halts on all inputs.

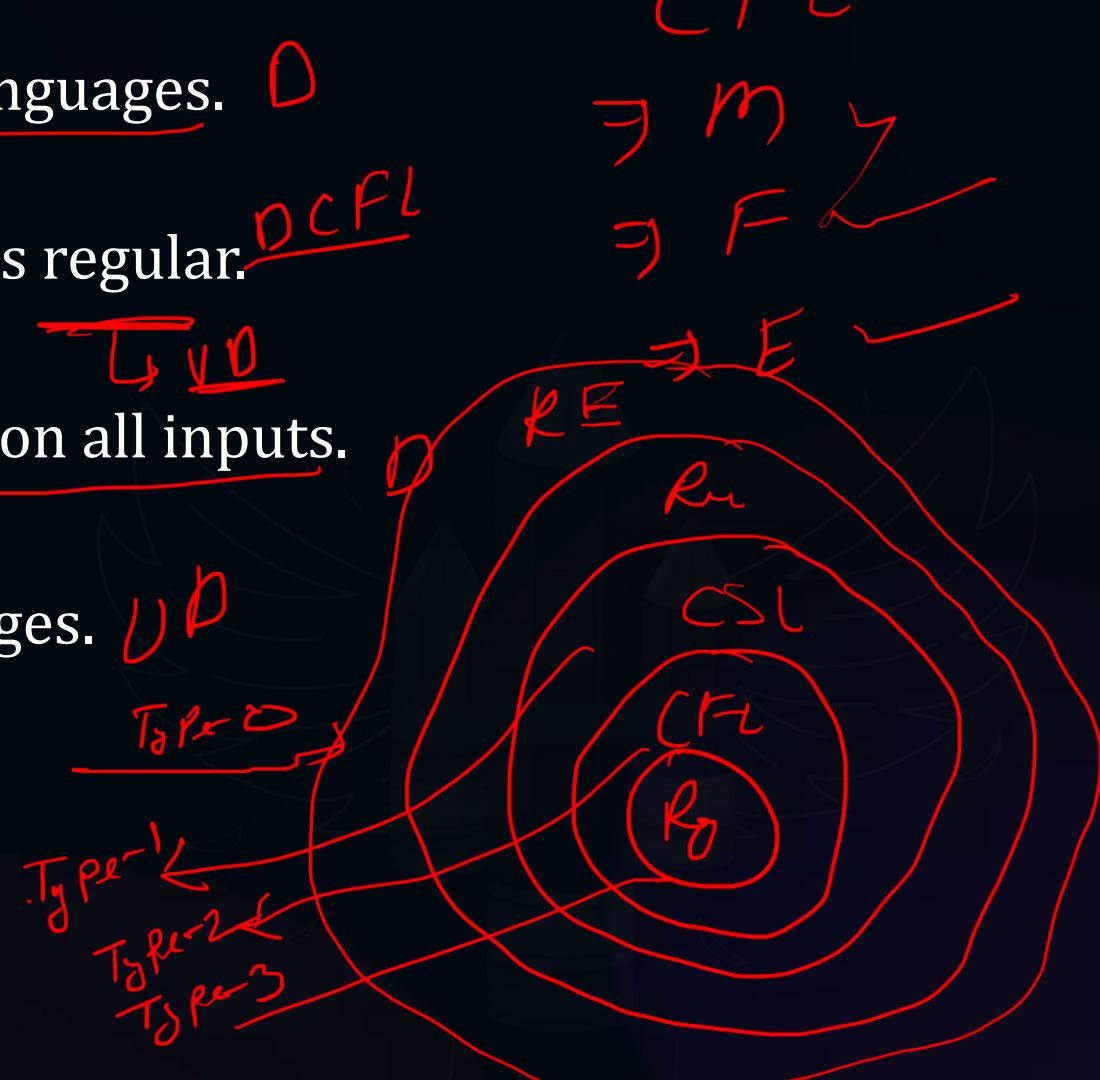
D

Membership problem for type-0 languages. UD

FSM

RF

Type 0
Type 1
Type 2
Type 3



[MCQ]



#Q. Which one of the following is UNDECIDABLE?

(B)

A

Decidable

Given a Turing machine M, a strings 's' and an integer k, M accepts s within k steps

B

~~Equivalence of two given Turing machines~~

UD

C

Language accepted by a given finite state machine is not empty

Decidable

D

Language generated by a context free grammar is non empty

Decidable

#Q. Consider the following decision problems:

- (P₁) Does a given finite state machine accept a given string? ✓
- (P₂) Does a given context free grammar generate an infinite number of strings

Which of the following statements is true?

A

Both (P₁) and (P₂) are decidable

C

Only (P₁) is decidable

B

Neither (P₁) nor (P₂) are decidable

D

Only (P₂) is decidable

[MCQ]



#Q. Which of the following problems is undecidable?

A

Membership problem for CFGs.

Dec

B

Ambiguity problem for CFGs.

UD ✓

C

Finiteness problem for FSAs.

Dec

D

Equivalence problem for FSAs.

Dec

CFL

F

M

E

R_n

✓

✓

✓

CSL

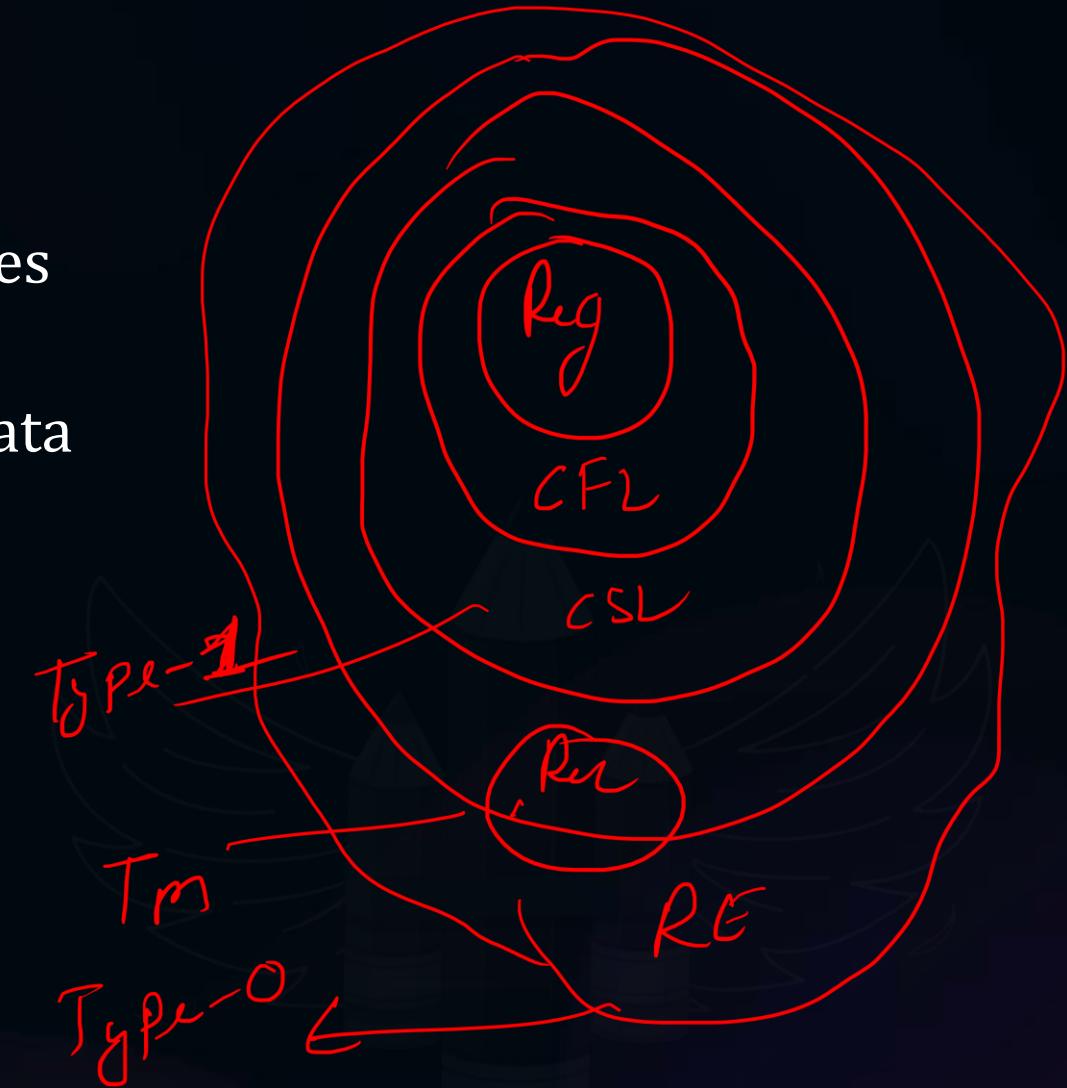
X

✓

X

#Q. Recursive languages are:

- A ~~a proper superset of context free languages~~
- B ~~always recognized by a pushdown automata~~
- C ~~also called type-0 languages~~
- D ~~recognizable by Turing machines~~



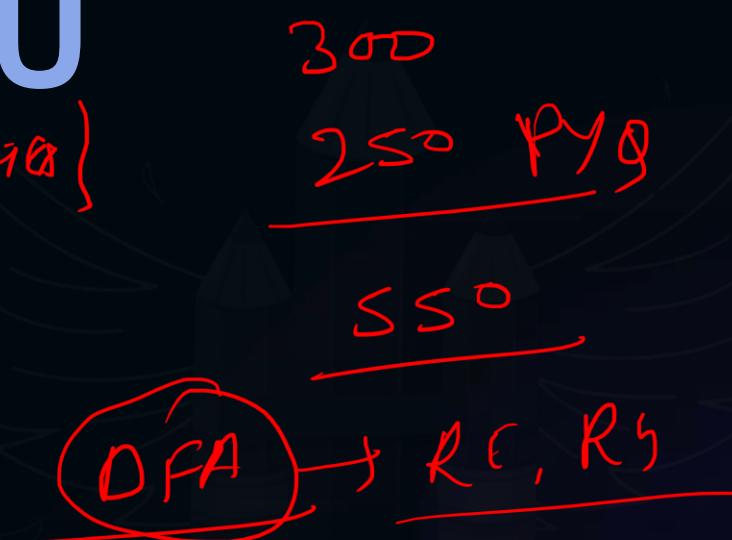
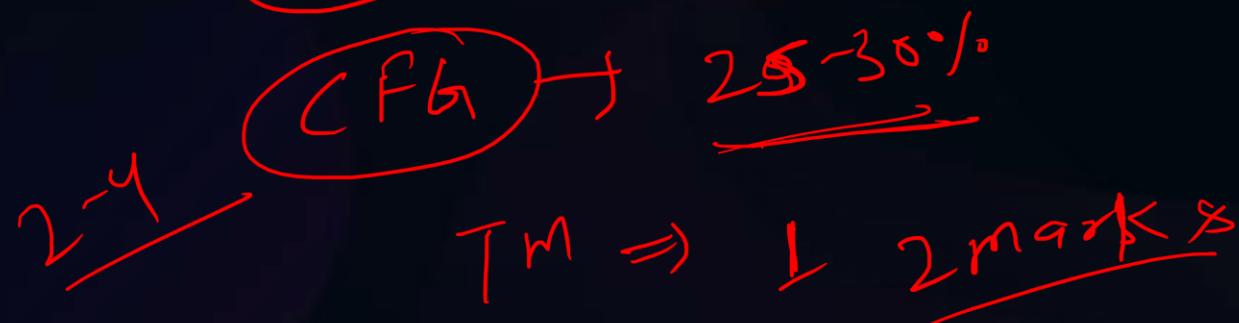
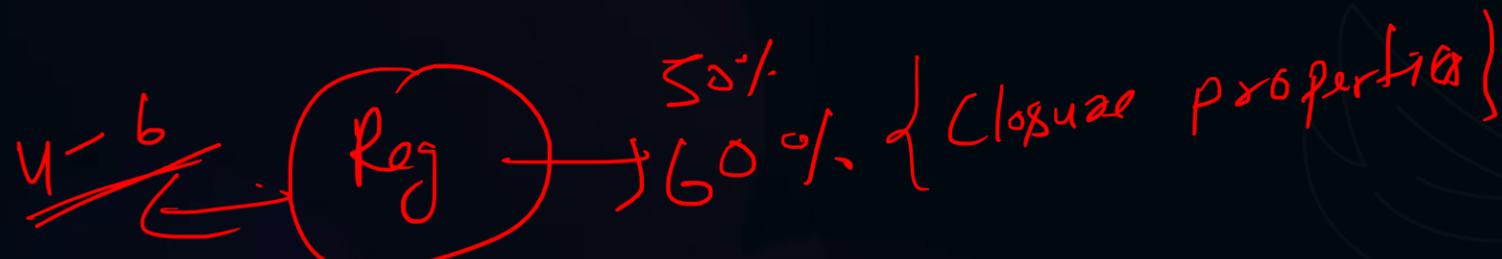
Classes over

$$\text{SWT} \Rightarrow \begin{cases} 50-60 \text{ problems} \\ 40-50 \text{ problems} \end{cases} \quad \text{CN} \Rightarrow \underline{200 \text{ problems}}$$

Telegram

100 problem

THANK - YOU



$\text{TM} \Rightarrow 1 \text{ 2 marks}$