

CS & IT ENGINEERING

THEORY OF COMPUTATION

Grammar

Lecture - 02



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Recap of Previous Lecture



Topic

?????

- ① Grammar
- ② Grammar Construction
- ③ Grammar → language

Topics to be Covered



Topic

Type of Grammar

Topic

?? Ambiguous Grammar

Topic

?? Simplification of Grammar

Topic

?? Normal form of Grammar.



Topic : Grammar

- ✓ Set of rules used to describe strings of a language is known as grammar.
- Formal definition of grammar is

$$G = (N, T, P, S)$$

- **N** :- non terminals (or) variables
- **T** :- Terminals
- **P** :- no. of productions
- **S** :- Starting symbol



Topic : Derivation

- The process of deriving strings from the given grammar known as derivation.
- The derivation can be either left most derivation or right most derivation
- **Left most derivation:**
It is the derivation in which left most non terminal is replaced by its R.H.S part at every step.
- **Right most derivation:**
It is a derivation in which right most non terminal is replaced by its R.H.S part at every step.

Derivation Tree (or)Parse Tree

- Tree representation of the derivation is known as derivation tree.
- All leaf node of the parse tree is known as yield of parse tree .
- while reading yield from left to right sentence of the grammar can be generate.

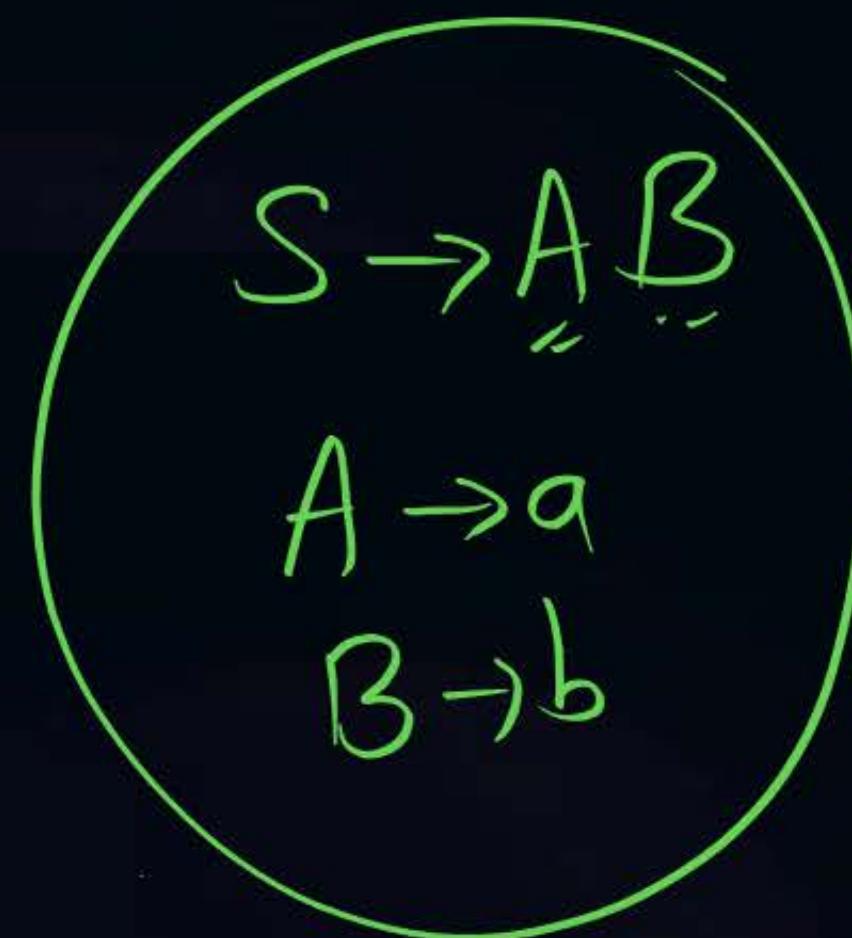
✓ → Regular Language → Regular Grammar

- ① ➤ Context free language there exist a grammar known as context free grammar.
- ② ➤ Context sensitive language there exist a grammar known as context sensitive grammar.
- ✓ ➤ For recursive enumerable language there exist a grammar known as unrestricted grammar.

#Q. Identify language generated by following grammar.

[NAT]

#Q. Construct grammar for the following languages.



$$\underline{L = \{aab\}}$$

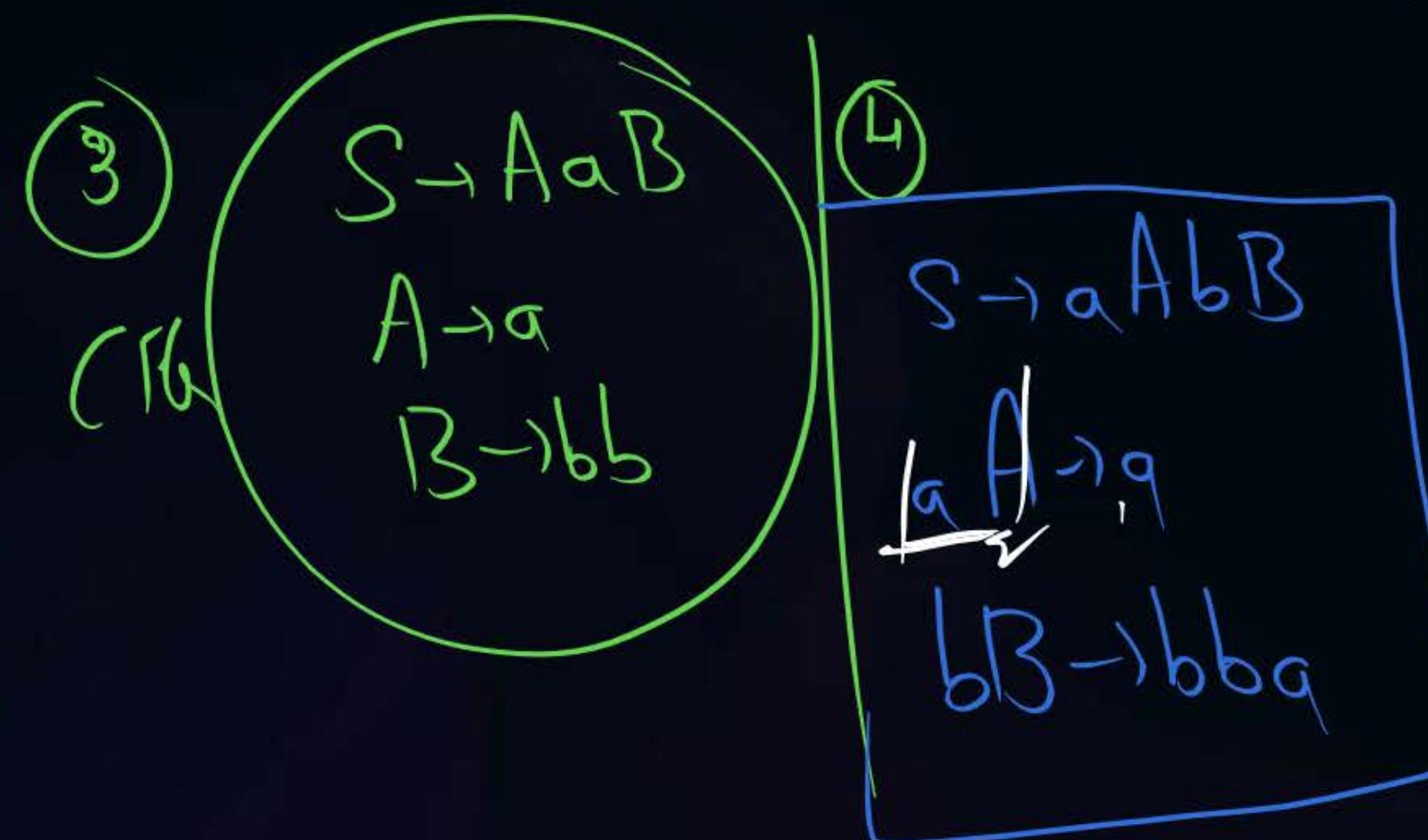
~~x not regular~~

x ③ Regular ?

② CFG ?

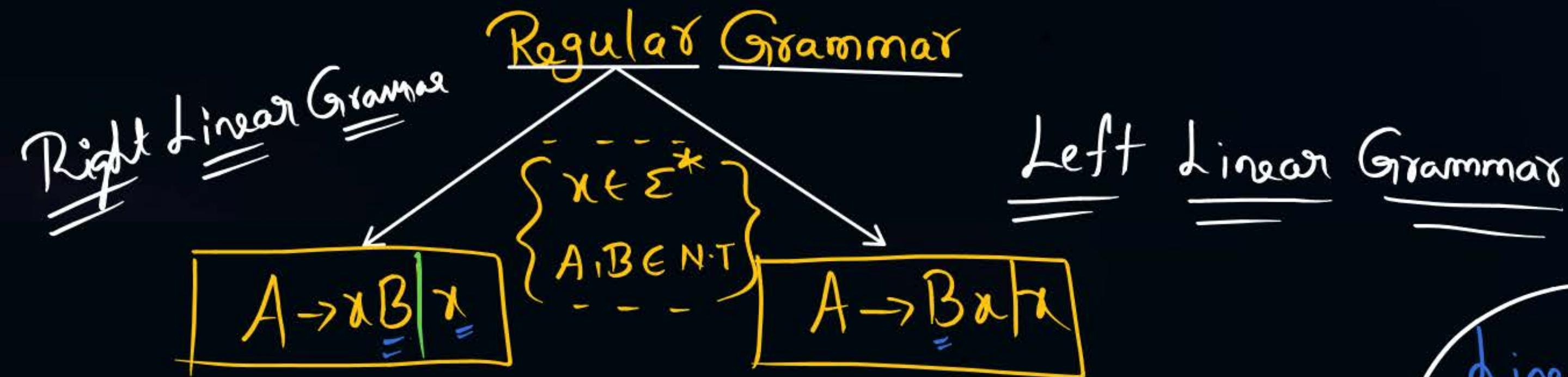
② $S \rightarrow aS \quad \overline{bS} \quad a$ $\stackrel{R.L.G}{\equiv}$

$\times a \text{ Reg}$
 $\times b \text{ CFB}$
 $\times c \text{ CSG}$
 $\checkmark d \text{ UFG}$



Types of Grammar

Type	Language (Grammars)	Form of Productions	Accepting Automata	Lang
③	Regular Gram	L.L(G) $A \rightarrow XBX^R$ $A \rightarrow BX^R$	Finite Automaton	Rg La
②	Context-free Grammar	$A \rightarrow \alpha$	$\alpha \in (V \cup T)^*$ Pushdown Automaton	CFL
①	Context-sensitive Grammar	$\alpha \Rightarrow^* \beta$ $ \alpha \leq \beta $	$\alpha, \beta \in (V \cup T)^*$ LBA	CSL
0 (R.E)	Unrestricted Gram	$\alpha \rightarrow \beta$	$\alpha, \beta \in (V \cup T)^*$ Turing machine	R.E
Type 0: $S \rightarrow AaBb$ $Aa \rightarrow aa$ $Bb \rightarrow b$		Type 1 (C.S.G) $S \rightarrow aAbB$ $aA \rightarrow aa$ $bB \rightarrow aba$	Type 2 (CFG) $S \rightarrow AaABA$ $A \rightarrow aBab$ $B \rightarrow ba$	Type 3



① $S \rightarrow aS \mid b$

② $S \rightarrow S a \mid b$

③ $S \rightarrow aS \mid Sb \mid a$

{ L: Linear but not regular }

④ $S \rightarrow Sa \mid Sb \mid \epsilon$

⑤ $S \rightarrow aSb \mid ab$

{ Linear but not regular }



{Chomsky Hierarchy}

Type 0 - Recursively enumerable

Unrestricted Gram.

Type 1 - Context-sensitive

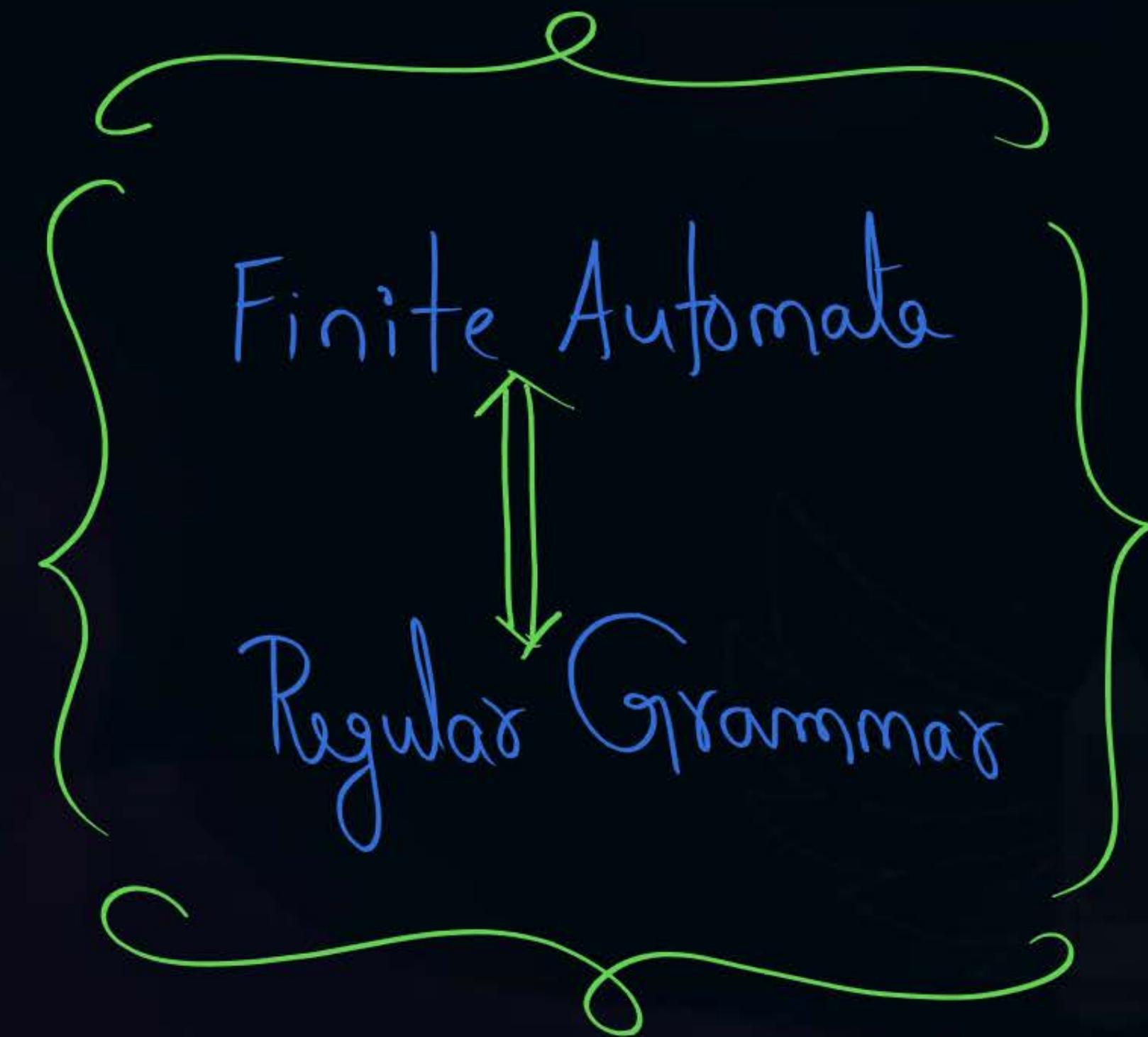
Gram.

Type 2 - Context-free

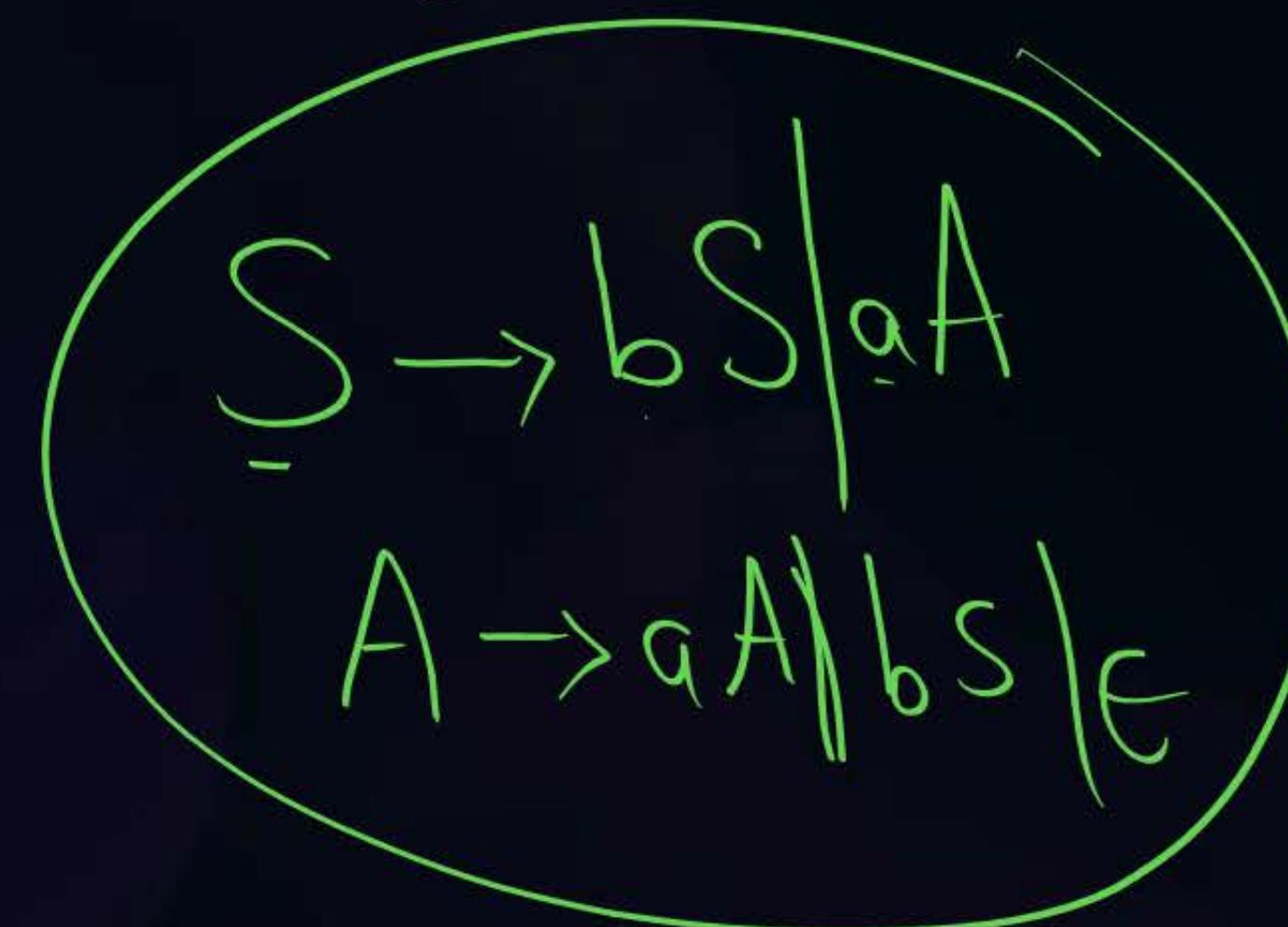
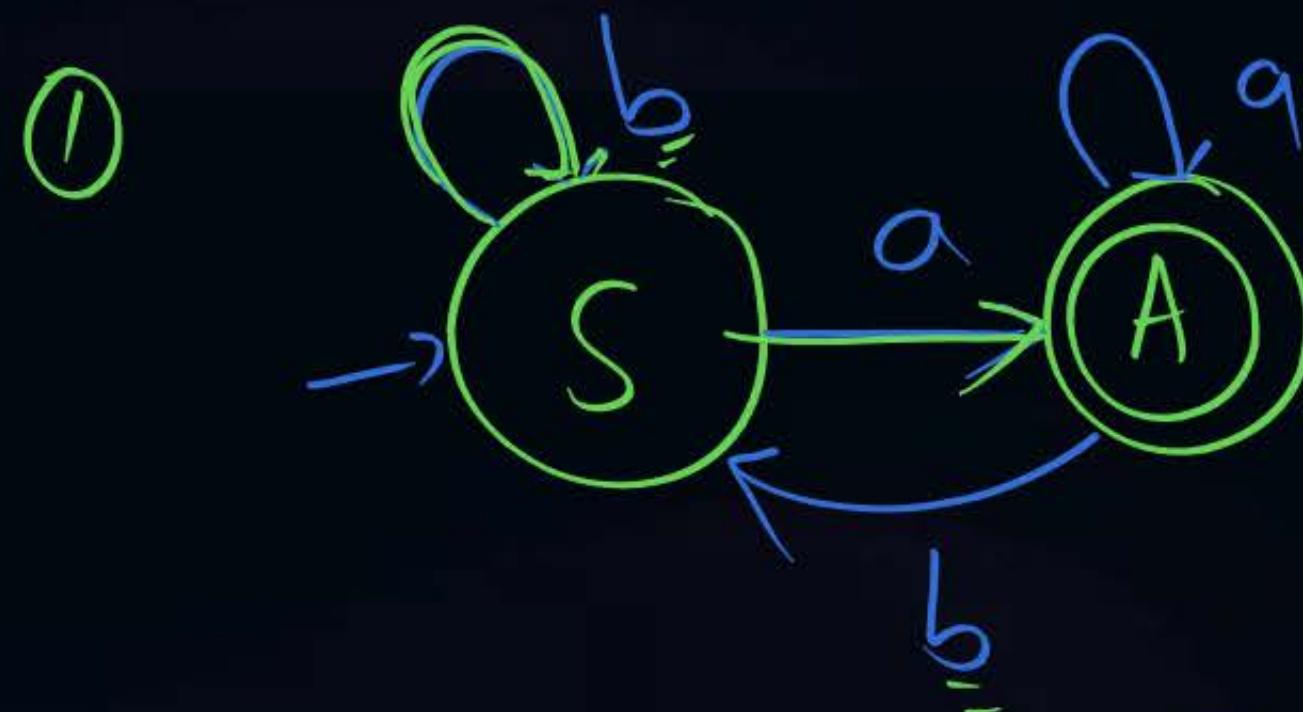
Gram.

Type 3 - Regular

Gram.



(Q) Construct Regular Grammar for given F.A



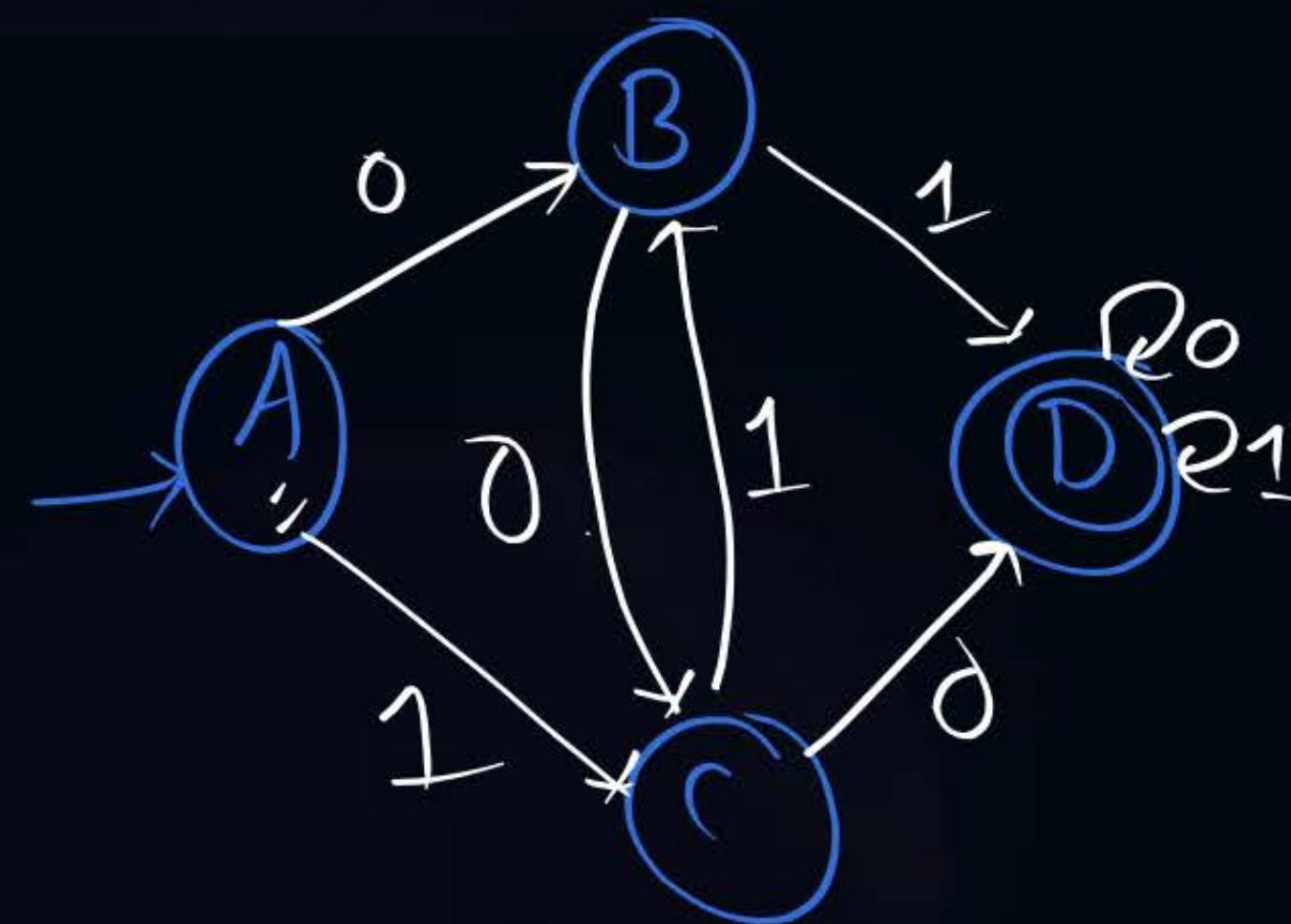
Grammar
(N, T, P, S)

N · T => {state}
T => Σ

P = transition
S = Initial state

(Q) Construct Regular Grammar for given F.A

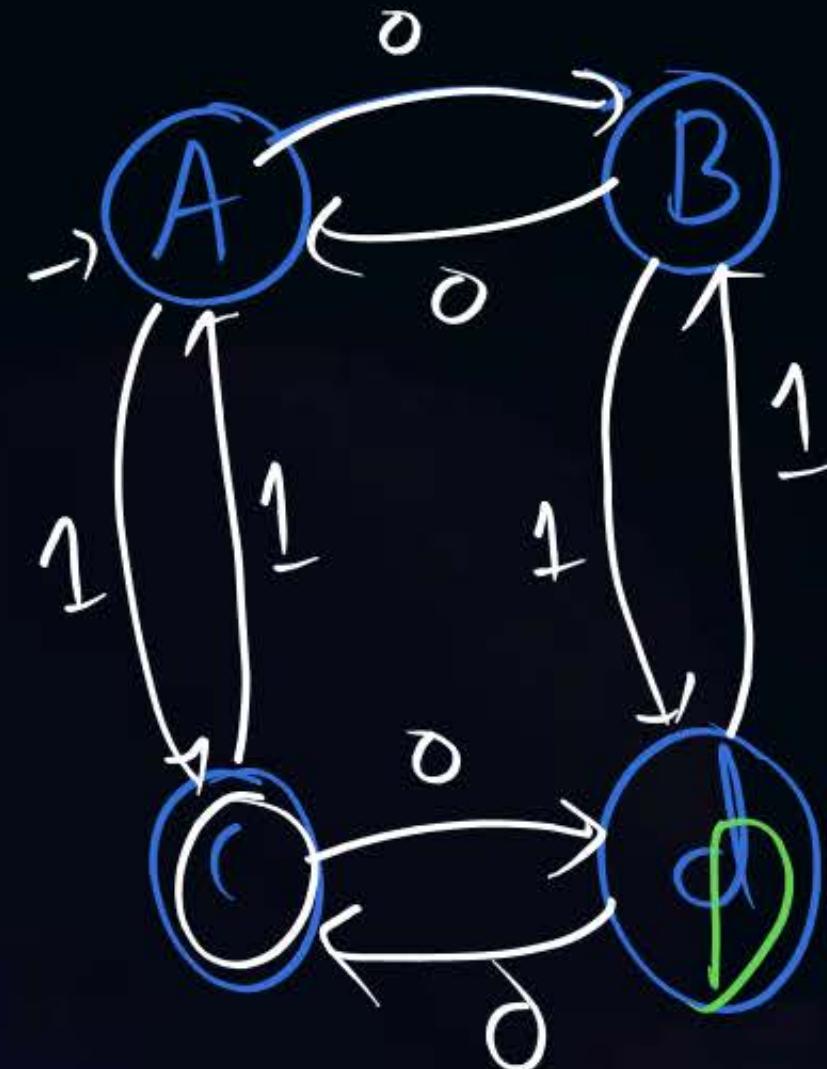
②



$$\Rightarrow \left\{ \begin{array}{l} A \rightarrow 0B|1C \\ B \rightarrow 0C|1D \\ C \rightarrow 0D|1B \\ D \rightarrow 0D|1D|\epsilon \end{array} \right. \quad \text{q producing}$$

(Q) Construct Regular Grammar for given F.A

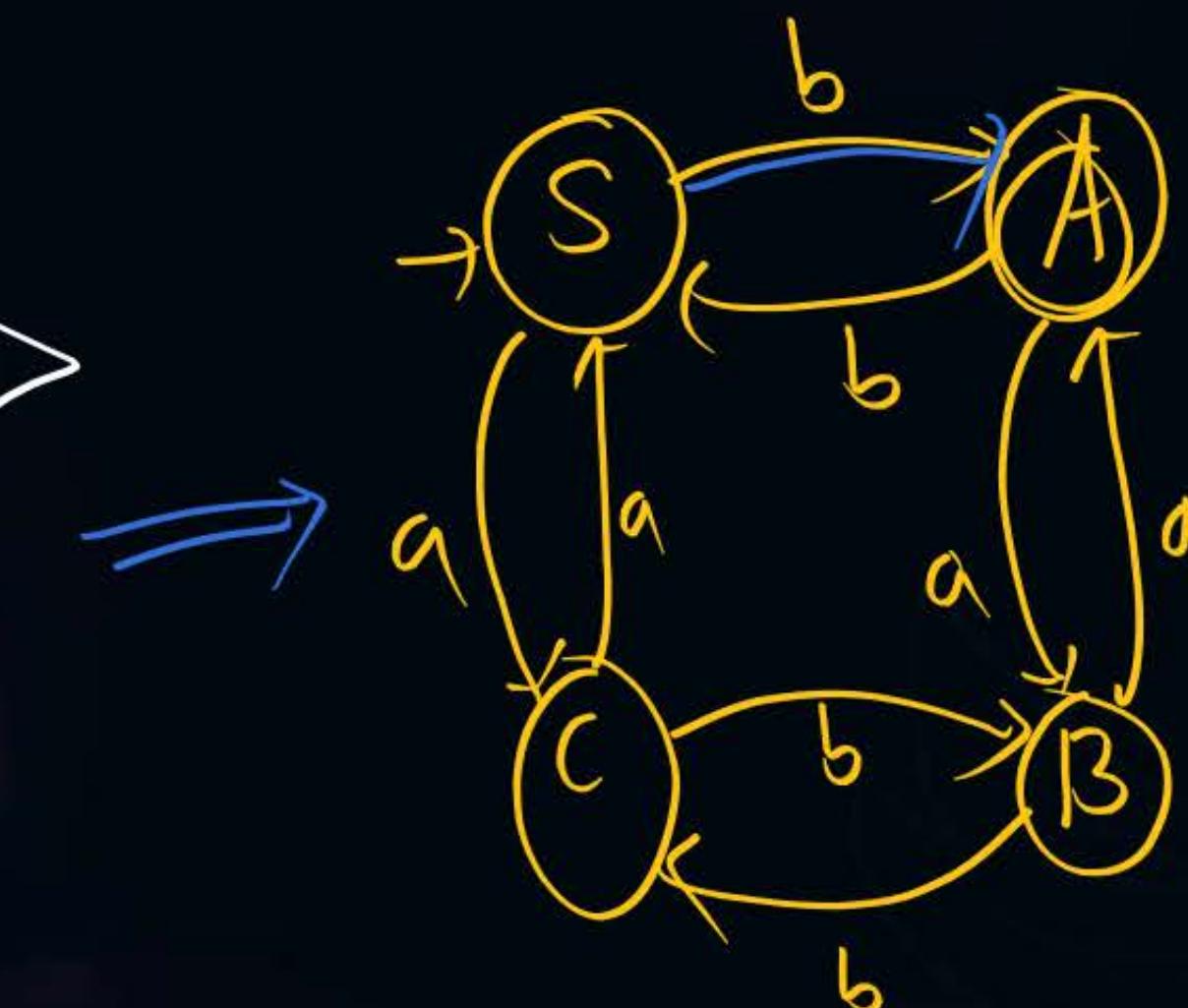
③



$$\left. \begin{array}{l} A \rightarrow \underline{0} \underline{B} \mid \underline{1} \underline{C} \\ B \rightarrow 0 \underline{A} \mid 1 \underline{D} \\ C \rightarrow 0 \underline{D} \mid 1 \underline{A} \mid \epsilon \\ D \rightarrow 0 \underline{C} \mid 1 \underline{B} \end{array} \right\} \text{q produk} =$$

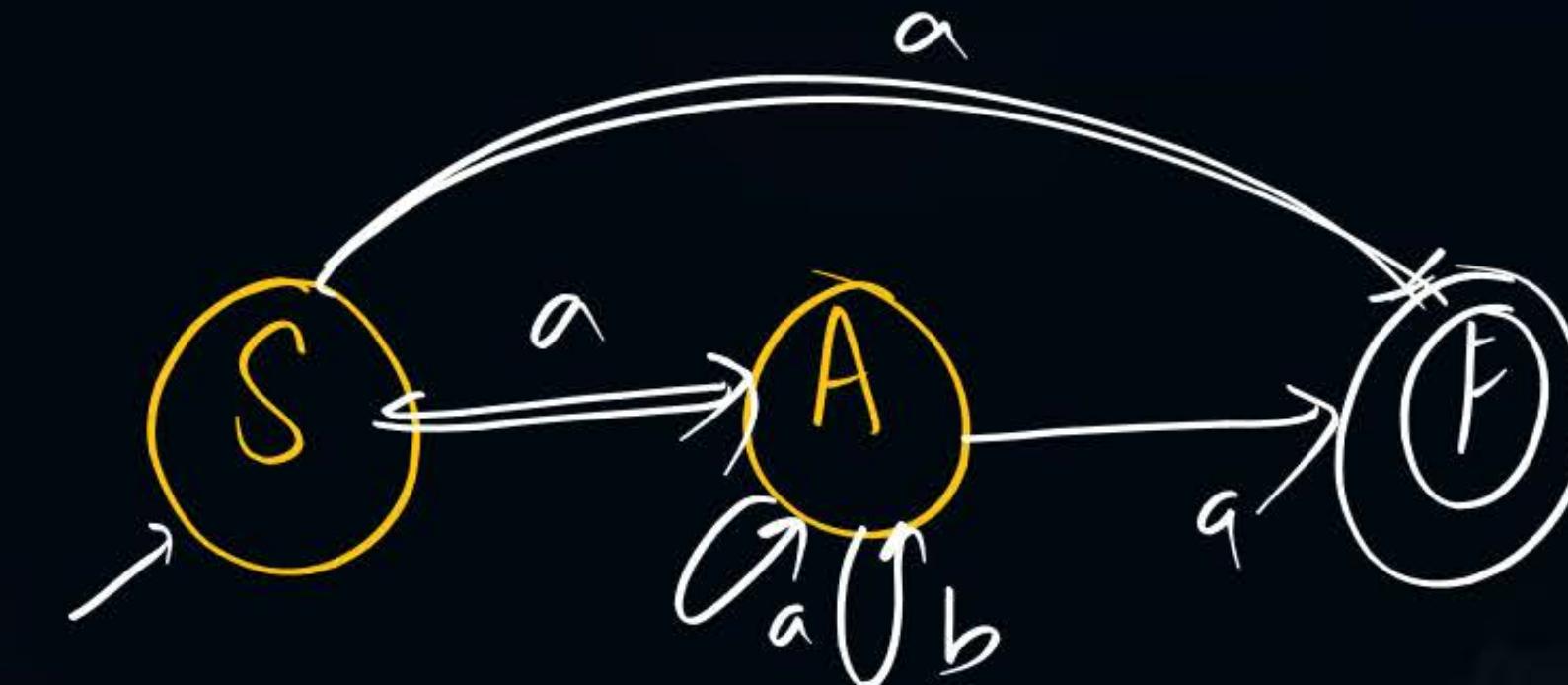
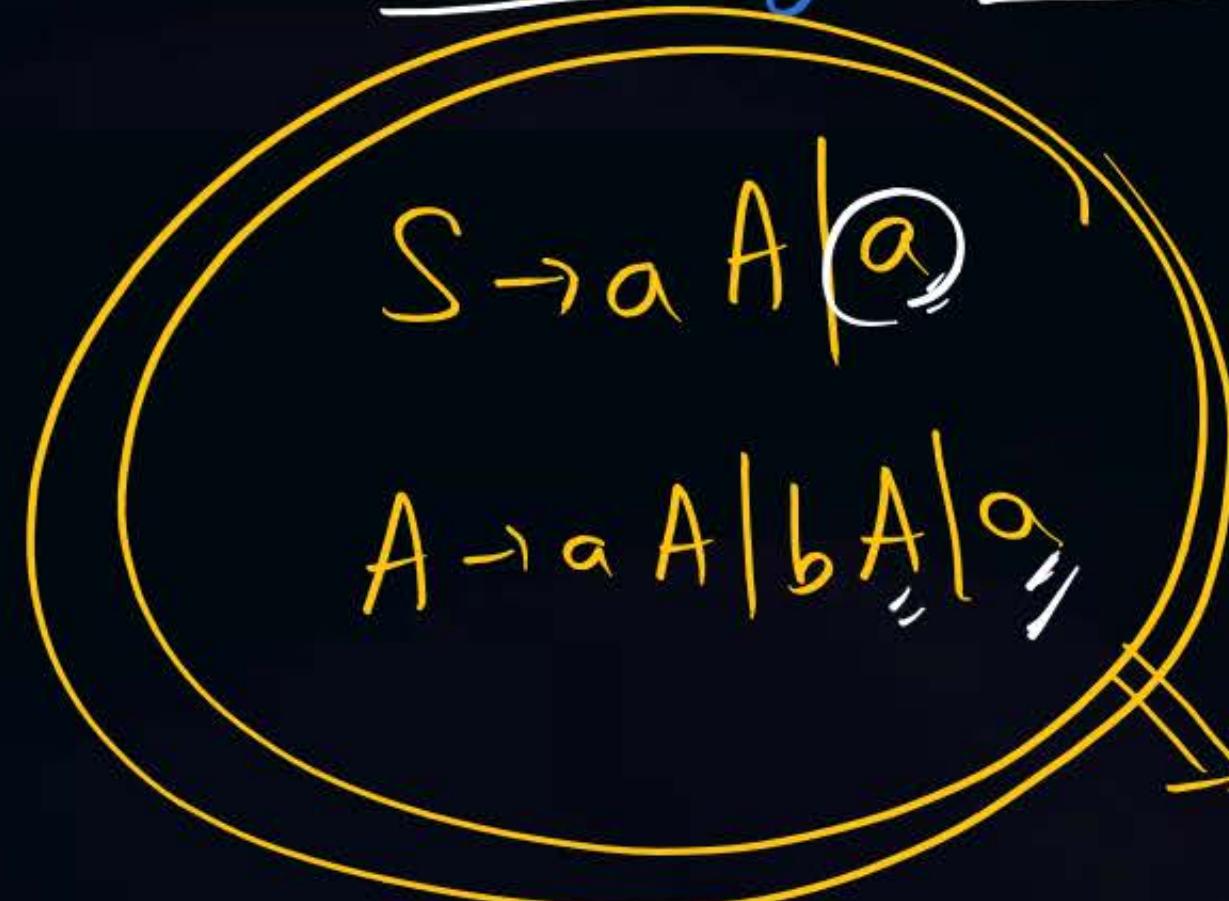
(Q) Identify language of given Grammar

$$\left. \begin{array}{l} S \rightarrow bA \mid aC \\ A \rightarrow bS \mid aB \mid \epsilon \\ B \rightarrow aA \mid bC \\ C \rightarrow bB \mid aS \end{array} \right\} \quad \epsilon$$



Has even (and) # $b \}$ odd a

(Q) Identify Regular Expression from given F.A?



f.a

$$a(a+b)^*a + a$$

Ambiguous Grammar

A Grammar is Said to be ambiguous for atleast one string if there exist more than one

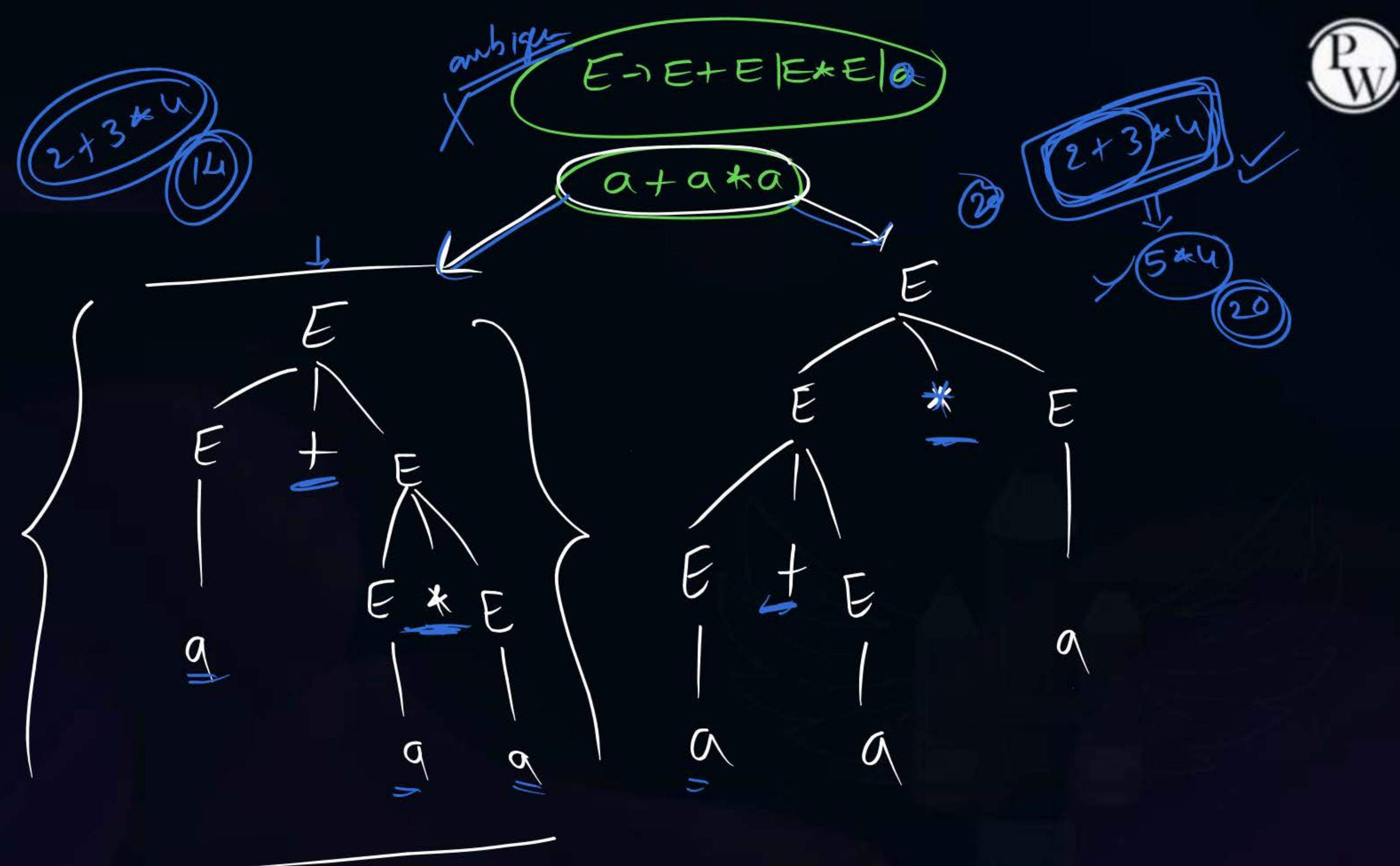
$\rightarrow \underline{L \cdot M \cdot D}$

(or)

$\rightarrow \underline{R \cdot M \cdot I}$

(or)

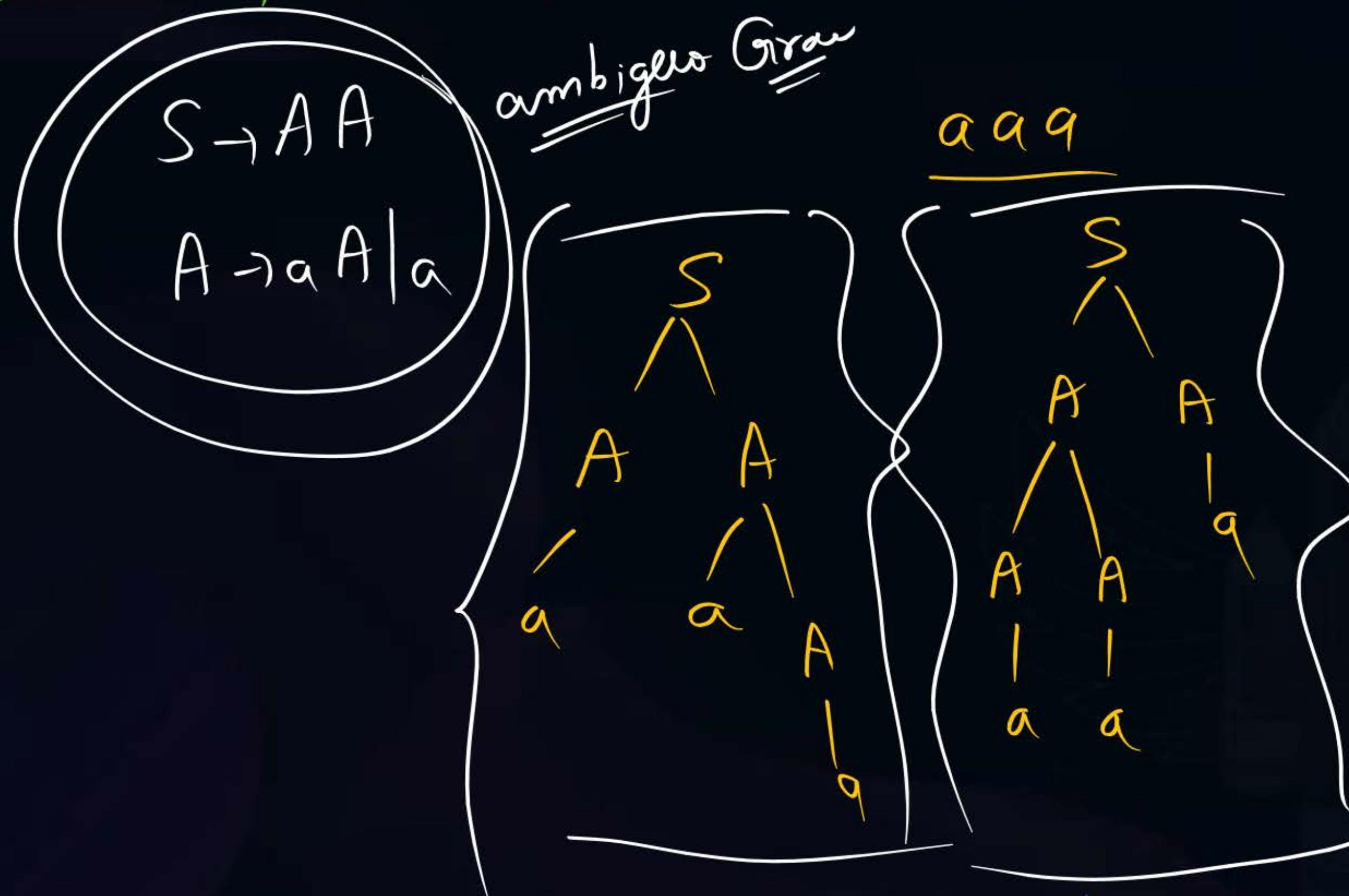
Parse tree



(Q) Verify given grammar if ambiguous (a) not?



(Q) Verify given grammar if ambiguous (or) not?



(Q) Verify given grammar if ambiguous (or) not?

$$\left. \begin{array}{l} S \rightarrow AaAb \mid BbBa \\ A \rightarrow \epsilon \\ B \rightarrow \epsilon \end{array} \right\}$$

Home work
=====

NOTE

Verifying given grammar ambiguous (a) not by

Undecidable Problem. No Algorithm exist

(Q) Construct Regular Grammar for given F.A



2 mins Summary



Topic One

Topic Two

Topic Three

Topic Four

Topic Five

Type

F.A.C.G

Ambiguity



THANK - YOU