

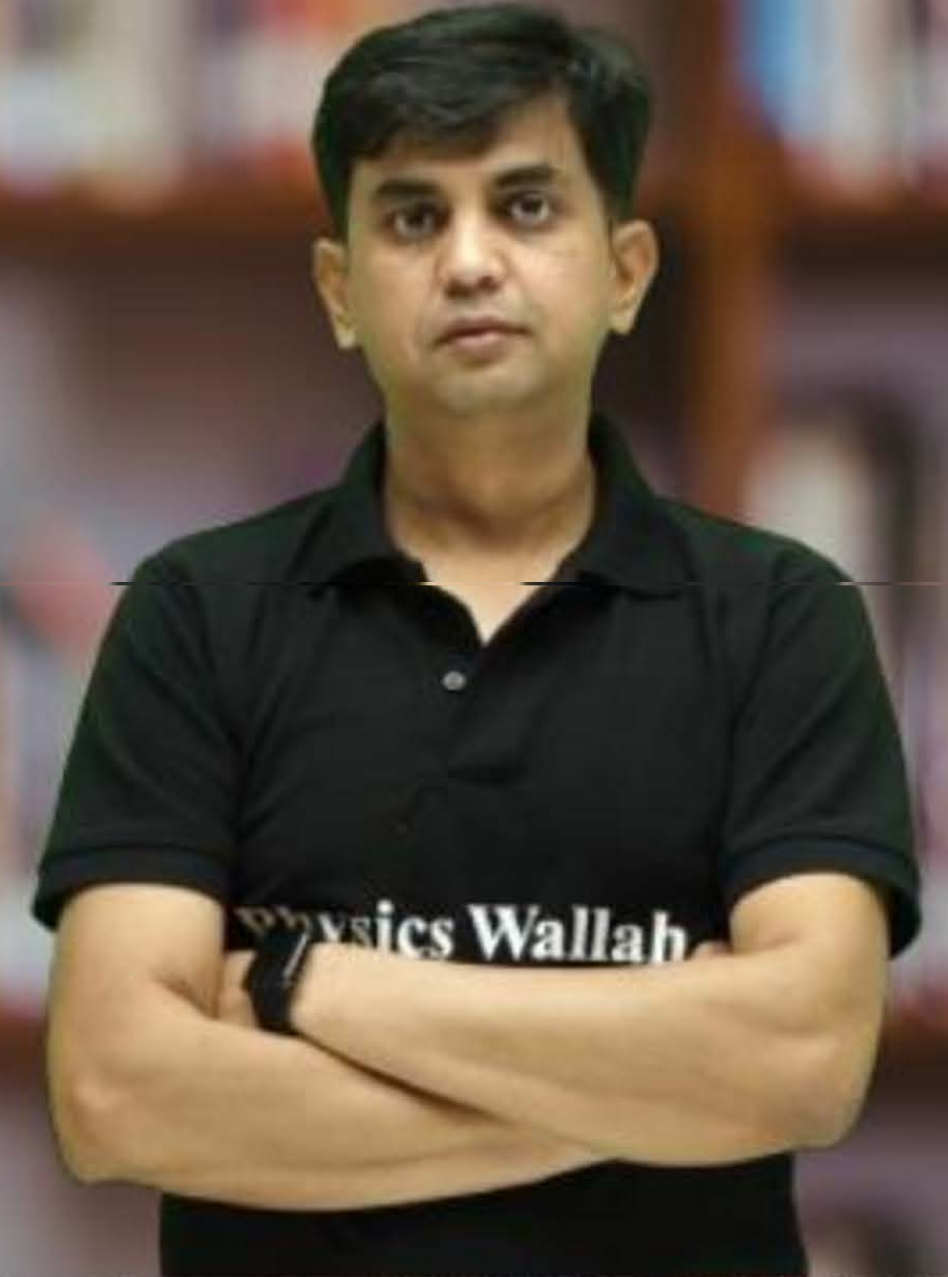
# CS & IT ENGINEERING



## Computer Network

### Flow Control

**DPP - 01 Discussion Notes**



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## [MCQ]



#Q. Consider two hosts are connected directly using point-to-point link. The link bandwidth is 10 Kbps, distance is 5 km, signal propagation speed is  $10^5$  meter per second and packet size is 200 bits than calculate transmission delay for a packet and one-way propagation delay (in millisecond) ?

☒ **A** 20 and 50

☐ **B** 20 and 500

☐ **C** 200 and 50

☐ **D** 200 and 500

$$t_x = \frac{200 \text{ bits}}{10 \text{ Kbps}} = \frac{200 \text{ bits}}{10 \times 10^3 \text{ bits/sec}} = 20 \text{ ms}$$

$$t_p = \frac{5 \text{ KM}}{10^5 \text{ m/sec}} = \frac{5 \times 10^3 \text{ m}}{10^2 \times 10^3 \text{ m/sec}} = 50 \text{ ms}$$

Ans: A



[NAT]



#Q. Consider two hosts S and D are connected directly using point-to-point link. The link bandwidth is 20 Kbps, signal propagation speed is  $10^6$  meter per second and distance is 40 km. Suppose host S wants to send a data of size 500 bits to host D in one packet. Calculate how much time (in millisecond) it takes to deliver the data completely from S to D?

$$t_x = \frac{500 \text{ bits}}{20 \text{ Kbps}} = \frac{500 \text{ bits}}{20 \times 10^3 \text{ bits/sec}} = 25 \text{ ms}$$

$$t_p = \frac{40 \text{ KM}}{10^6 \text{ m/sec}} = \frac{40 \times 10^3 \text{ m}}{10^3 \times 10^3 \text{ m/sec}} = 40 \text{ ms}$$

$$\text{Total time} = (t_x + t_p) = (25 + 40) \text{ ms} = 65 \text{ ms}$$

Ans = 65



[NAT]



#Q. Consider two hosts S and D are connected via one packet switch (takes negligible processing delay). Each link bandwidth is 1 Mbps, signal propagation speed is  $4 * 10^6$  meter per second and length of each link is 20 km. Suppose host S wants to send a file of size 8000 bytes to host D, by dividing the file into 2000 bits packets. Calculate how much time (in millisecond) it takes to deliver the file completely from S to D?

$$t_x = \frac{2000 \text{ bits}}{1 \text{ Mbps}} = \frac{2 * 10^3 \text{ bits}}{10^6 \text{ bits/sec}} = 2 \text{ ms}$$

$$t_p = \frac{20 \text{ km}}{4 * 10^6 \text{ m/sec}} = \frac{20 * 10^3 \text{ m}}{4 * 10^6 \text{ m/sec}} = 5 \text{ ms}$$

$$N = \frac{8000 \text{ bytes}}{2000 \text{ bits}} = \frac{64 * 10^3 \text{ bits}}{2 * 10^3 \text{ bits}} = 32$$

$$\begin{aligned} \text{Total Time} &= (N * t_x + t_p) + (t_x + t_p) \\ &= (32 * 2 + 5) + (2 + 5) \text{ ms} \\ &= 76 \text{ ms} \end{aligned}$$

$$\boxed{\text{Ans} = 76}$$



[NAT]



#Q. Consider two hosts S and D are connected via two routers over a single path, routers takes negligible processing delay. Each link bandwidth is 100 Kbps, signal propagation speed is  $10^6$  meter per second and length of each link is 10 km. Suppose host S wants to send a file of size 5000 bytes to host D, by dividing the file into 1000 bits packets. Calculate how much time (in millisecond) it takes to deliver the file completely from S to D?

$$t_x = \frac{1000 \text{ bits}}{100 \text{ Kbps}} = \frac{1000 \text{ bits}}{100 \times 10^3 \text{ bits/sec}} = 10 \text{ ms}$$

$$t_p = \frac{10 \text{ KM}}{10^6 \text{ m/sec}} = \frac{10 \times 10^3 \text{ m}}{10^3 \times 10^3 \text{ m/sec}} = 10 \text{ ms}$$

$$N = \frac{5000 \text{ bytes}}{1000 \text{ bits}} = \frac{40 \times 10^3 \text{ bits}}{10^3 \text{ bits}} = 40$$

$$\begin{aligned} \text{Total time} &= (N \times t_x + t_p) + (t_x + t_p) + (t_x + t_p) \\ &= (40 \times 10 + 10) + (10 + 10) + (10 + 10) \text{ ms} \\ &= 450 \text{ ms} \end{aligned}$$

Ans = 450



[NAT]



#Q. Consider two hosts are connected via point-to-point link using Stop-and-Wait protocol for flow control. The transmission rate of the link is 5 Kbps and the one-way propagation delay of the link is 50 milliseconds. The frame size is 500 bits. Calculate channel utilization in percentage (rounded off to nearest integer) ?

$$t_x = \frac{500 \text{ bits}}{5 \text{ Kbps}} = \frac{500 \text{ bits}}{5 \times 10^3 \text{ bits/sec}} = 100 \text{ ms}$$

$$t_p = 50 \text{ ms}$$

$$\begin{aligned} \text{Cycle time} &= [t_x + 2t_p] \\ &= (100 + 2 \times 50) \text{ ms} \\ &= 200 \text{ ms} \end{aligned}$$

$$\begin{aligned} \eta &= \frac{t_x}{\text{cycle time}} \\ &= \frac{100 \text{ ms}}{200 \text{ ms}} = \frac{1}{2} \\ &= \frac{1}{2} \times 100\% \\ &= 50\% \end{aligned}$$

Ans = 50



[NAT]



#Q. Consider two hosts are connected via point-to-point link using Stop-and-Wait protocol for flow control. The transmission rate of the link is 2 Mbps and frame size is 4000 bytes. To achieve 50 percent link utilization, calculate one-way propagation delay in milliseconds?

$$t_x = \frac{4000 \text{ bytes}}{2 \text{ Mbps}} = \frac{32 \times 10^3 \text{ bits}}{2 \times 10^6 \text{ bits/sec}} = 16 \text{ ms}$$

$$\eta = \frac{t_x}{\text{cycle time}}$$

$$\frac{1}{2} = \frac{t_x}{(t_x + 2t_p)}$$

$$t_x = 2t_p$$

$$t_p = \frac{t_x}{2}$$

$$t_p = \frac{16 \text{ ms}}{2} = 8 \text{ ms}$$

Ans = 8



[NAT]



#Q. Consider two hosts are connected via point-to-point link using Stop-and-Wait protocol for flow control. The transmission delay for a frame and one-way propagation delay of the link are 15 and 30 milliseconds respectively. If link bandwidth is 50 Kbps then calculate throughput of the link in Kbps?

$$\begin{array}{l|l} t_x = 15 \text{ ms} & \text{cycle time} = (t_x + 2t_p) = (15 + 2 \times 30) \text{ ms} \\ t_p = 30 \text{ ms} & = 75 \text{ ms} \end{array}$$

$$\eta = \frac{t_x}{\text{cycle time}} = \frac{15 \text{ ms}}{75 \text{ ms}} = \frac{1}{5}$$

$$\begin{aligned} \text{Throughput} &= \eta * \text{Bandwidth} \\ &= \frac{1}{5} * 50 \text{ Kbps} \\ &= 10 \text{ Kbps} \end{aligned}$$

Ans = 10



[NAT]



#Q. Consider two hosts are connected via point-to-point link using Stop-and-Wait protocol for flow control. The link bandwidth is 4 Kbps, one-way propagation delay of the link are 800 milliseconds. To achieve 50 percent link utilization calculate packet size in bytes?

$$\eta = \frac{t_x}{\text{cycle time}}$$

$$\frac{1}{2} = \frac{t_x}{(t_x + 2t_p)}$$

$$t_x = 2t_p$$

$$\text{Packet size} = (2t_p) * \text{Bandwidth}$$

$$= (2 * 800 \text{ ms}) * 4 \text{ kbps}$$

$$= 64 * 10^2 * 10^{-3} \text{ sec} * 10^3 \text{ bits/sec}$$

$$= 6400 \text{ bits}$$

$$= 800 \text{ bytes}$$

Ans = 800



[NAT]



#Q. Consider a packet error probability is 0.3 in Stop-and-Wait ARQ. What should be minimum number of transmission required to transmit 140 packets?

$$\begin{aligned} P_b &= 0.3 \\ N &= 140 \end{aligned} \quad \left| \begin{aligned} \text{Total transmissions} &= \\ &= N + N * P_b + N * P_b^2 + N * P_b^3 + \dots \\ &= N * [1 + P_b + P_b^2 + P_b^3 + \dots] \\ &= N * \frac{1}{(1 - P_b)} = \frac{N}{(1 - P_b)} \\ &= \frac{140}{(1 - 0.3)} = \frac{140}{0.7} = 200 \end{aligned} \right.$$

Ans = 200



# [MCQ]



#Q. Consider two hosts are connected via point-to-point link using Sliding Window protocol for flow control. The transmission delay for a frame and propagation delay of the link are 10 and 30 milliseconds. Calculate optimal window size?

**A**

1

$$t_x = 10 \text{ ms}$$
$$t_p = 30 \text{ ms}$$

$$\text{Cycle time} = (t_x + 2t_p) = (10 + 2 \times 30) \text{ ms}$$
$$= 70 \text{ ms}$$

**B**

3

$$\text{optimal window size} = \left\lceil \frac{\text{Cycle time}}{t_x} \right\rceil$$

**C**

4

$$= \left\lceil \frac{70 \text{ ms}}{10 \text{ ms}} \right\rceil = 7$$

**D**

7

Ans: D



[NAT]



#Q. Consider two hosts A and B are connected via point-to-point link using Go-Back-4-ARQ flow control strategy with transmit window size 4. Host A wants to send a file (divided into 10 packets) to host B, if every 6th packet that A transmits gets lost (but no ACKs from B ever get lost), then what is the number of packets that A will transmit for sending the file to B?

1	2	3	4	5	6	7	8	9
✓	✓	✓	✓	✓	X	Discard		

6	7	8	9	10
✓	✓	X	discard	
<hr/>				
		8	9	10
		✓	✓	✓

Ans = 17



## [MCQ]



#Q. Consider two hosts connected directly through point-to-point link uses Go Back N ARQ for flow control. The transmission delay and propagation delay (in millisecond) are 2 and 250. What should be minimum number of bits required for sequence number field to achieve full utilization?

**A**

7

$$t_x = 2 \text{ ms}, t_p = 250 \text{ ms}$$
$$\text{Cycle time} = (t_x + 2t_p)$$
$$= (2 + 2 \times 250) \text{ ms} = 502 \text{ ms}$$

**B**

8

**C**

9

**D**

10

$$\text{Optimal window size (N)}$$
$$= \left\lceil \frac{\text{Cycle time}}{t_x} \right\rceil = \left\lceil \frac{502 \text{ ms}}{2 \text{ ms}} \right\rceil$$

$$N = 251$$

$$\text{Min}^m \text{ no. of bits for sequence no. field}$$
$$= \lceil \log_2 (N+1) \rceil \text{ bits}$$
$$= \lceil \log_2 (252) \rceil \text{ bits}$$
$$= 8 \text{ bits}$$

Ans: B



[NAT]



#Q. Consider a 2 Mbps communication link with one-way propagation delay of 120 milliseconds. Go back N ARQ protocol is used on this link to send data with a frame size of 2000 bits. Neglect the transmission time of acknowledgment. The minimum number of bits required for the sequence number field to achieve 50% utilization is \_\_\_\_\_.

$$t_x = \frac{2000 \text{ bits}}{2 \text{ Mbps}} = \frac{2 \times 10^3 \text{ bits}}{2 \times 10^6 \text{ bits/sec}} = 1 \text{ ms} \quad | \quad t_p = 120 \text{ ms}$$

$$\text{Cycle time} = (t_x + 2t_p) = (1 + 2 \times 120) \text{ ms} = 241 \text{ ms}$$

For 50% utilization

$$N = \left\lceil \frac{\text{Cycle time}}{2t_x} \right\rceil = \left\lceil \frac{241 \text{ ms}}{2 \times 1 \text{ ms}} \right\rceil = 121$$

$$\begin{aligned} \text{Min}^m \text{ no of bits for seq. no} &= \lceil \log_2(N+1) \rceil \text{ bits} \\ &= \lceil \log_2(122) \rceil \text{ bits} \\ &= 7 \text{ bits} \end{aligned}$$

Ans = 7



## [MCQ]



#Q. Consider two hosts connected directly through point-to-point link uses Selective Repeat ARQ for flow control. The transmission delay and propagation delay (in millisecond) are 4 and 250. What should be minimum number of bits required for sequence number field to achieve full utilization?

**A**

5

$$t_x = 4 \text{ ms}, t_p = 250 \text{ ms}$$
$$\text{Cycle time} = (t_x + 2t_p)$$
$$= (4 + 2 \times 250) \text{ ms} = 504 \text{ ms}$$

**B**

6

**C**

7

**D**

8

Optimal window size

$$N = \left\lceil \frac{\text{Cycle time}}{t_x} \right\rceil = \left\lceil \frac{504 \text{ ms}}{4 \text{ ms}} \right\rceil$$

$$N = 126$$

$$\begin{aligned} \text{Min}^m \text{ no of bits for seq. no} \\ &= \lceil \log_2(2N) \rceil \text{ bits} \\ &= \lceil \log_2(252) \rceil \text{ bits} \\ &= 8 \text{ bits} \end{aligned}$$

Ans: D



[NAT]



#Q. Consider a 1 Mbps communication link with one-way propagation delay of 80 milliseconds. Selective repeat protocol is used on this link to send data with a frame size of 2000 bits. Neglect the transmission time of acknowledgment. The minimum number of bits required for the sequence number field to achieve 100% utilization is \_\_\_\_\_.

$$t_x = \frac{2000 \text{ bits}}{1 \text{ Mbps}} = \frac{2 \times 10^3 \text{ bits}}{10^6 \text{ bits/sec}} = 2 \text{ ms} \quad | \quad t_p = 80 \text{ ms}$$

$$\text{Cycle time} = (t_x + 2t_p) = (2 + 2 \times 80) \text{ ms} = 162 \text{ ms}$$

Optimal window size

$$N = \left\lceil \frac{\text{Cycle time}}{t_x} \right\rceil = \left\lceil \frac{162 \text{ ms}}{2 \text{ ms}} \right\rceil$$

$$N = 81$$

$$\begin{aligned} \text{Min}^m \text{ no. of bits for seq. no.} &= \lceil \log_2(2N) \rceil \text{ bits} \\ &= \lceil \log_2(162) \rceil = 8 \text{ bits} \end{aligned}$$

$$\boxed{\text{Ans} = 8}$$





**THANK - YOU**