

CS & IT ENGINEERING



Operating System

File System

(One Shot)

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Recap of Previous Lecture



Topic

Inverted Paging

Topic

Hashed Page table

Topics to be Covered



Topic

File System

Topic

Blocks and File Allocation Methods

Topic

Disk Scheduling

Topic

SSTF, Scan, Look



Topic : File



A file is a named collection of related information that is recorded on secondary storage.



Topic : File Attributes

1. Name
2. Extension
3. Size
4. Date
5. Author
6. Created, Modified, Accessed
7. Attributes: Read-only, hidden
8. Default Program
9. Security Details





Topic : File Directory



↗ folder

Collection of files



Topic : File System



Module of OS which manages, controls and organizes files and related structures



Topic : File System

The image shows a Windows File Explorer window with a dark theme. On the left, there's a sidebar with a 'Recent' section containing 'File Explorer', 'This PC', 'OneDrive', 'Recycle Bin', and 'Control Panel'. Below it are sections for 'Folders (7)' and 'Devices and drives (1)'. The main area displays the contents of the 'Desktop' folder, which includes '3D Objects', 'Downloads', 'Videos', 'Music', 'Documents', and 'Pictures'. It also shows a summary of the 'Windows-SSD (C:)' drive with 149 GB free of 237 GB. The background features a large, faint illustration of a leaf.

- ▼ Folders (7)
 - 3D Objects
 - Downloads
 - Videos
- Desktop
- Documents
- Music
- Pictures
- ▼ Devices and drives (1)
 - Windows-SSD (C:) 149 GB free of 237 GB





Topic : Types of File Systems

1. FAT32
2. NTFS → windows
3. HFS+ → iOS or macOS
4. Ext2 / Ext3 / Ext4 → linux / unix
5. Swap





Topic : File System



Windows-SSD (C:) Properties X

General		Tools	Hardware	Sharing
	Windows-SSD			
Type:	Local Disk			
File system:	NTFS 			
Used space:	93,87,61,70,752 bytes	87.4 GB		
Free space:	1,60,84,49,55,648 bytes	149 GB		
Capacity:	2,54,72,11,26,400 bytes	237 GB		
Drive C: Disk Cleanup				
<input type="checkbox"/> Compress this drive to save disk space				
<input checked="" type="checkbox"/> Allow files on this drive to have contents indexed in addition to file properties				

OK Cancel Apply



Topic : Disk Formatting



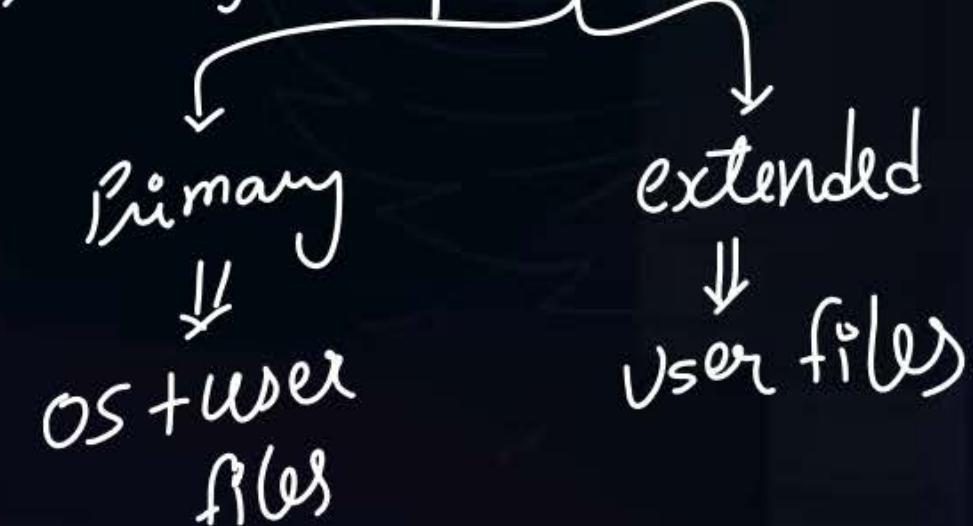
low level formatting
(physical)

creating tracks and sectors
on disk.

Done by manufacturer

high level formatting
(logical)

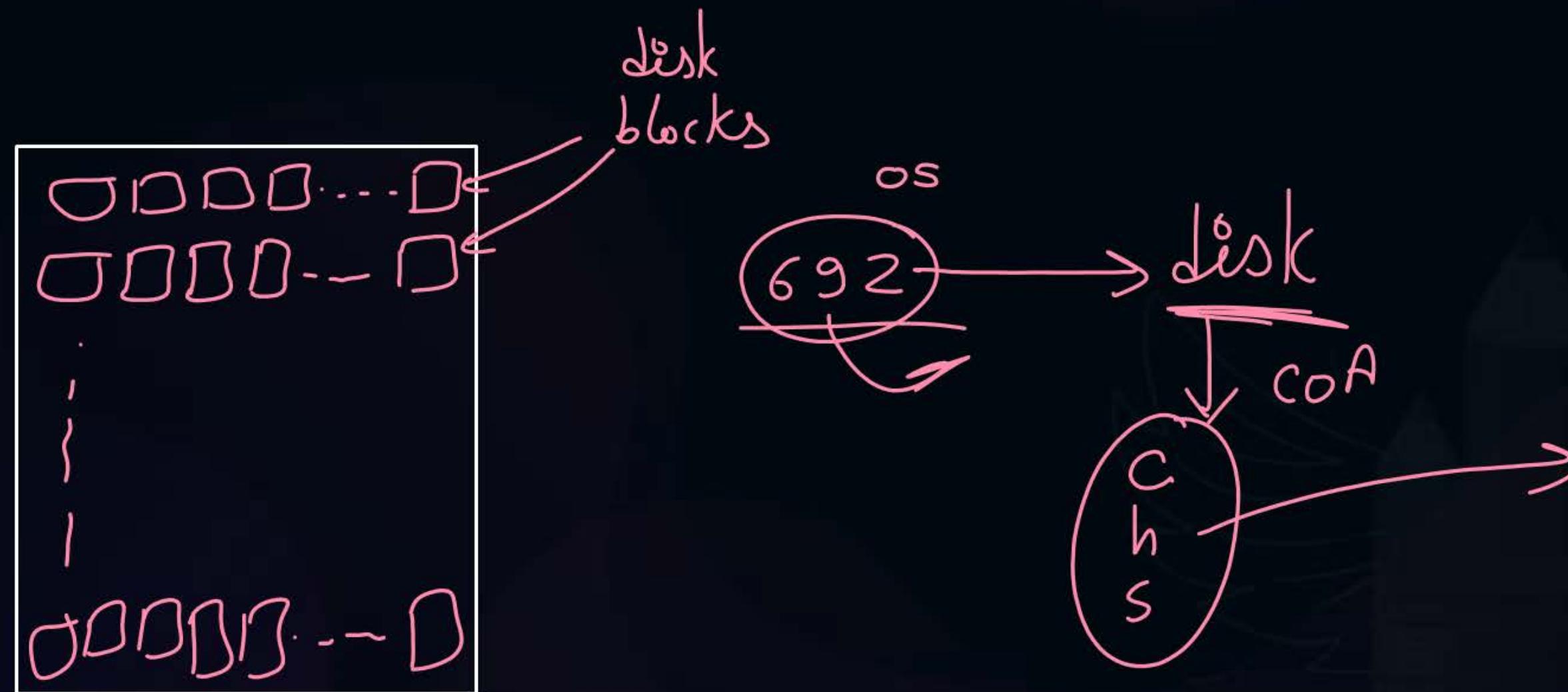
- ⇒ creating drives (partitions)
on disk like C:\ E:\ F:\
- creating blocks
- ⇒ 2 types of partitions





Topic : Disk Blocks

smallest logical unit of disk ($1 \text{ sector} \Rightarrow 1 \text{ block}$)





Topic : Disk Blocks

Number of disk blocks = 2^{16}

Size of each block = 1KB

Total Size of disk?

Disk block address (D.B.A.) = 16 bits

$$2^{16} * 1KB = 2^{26}B = 64MB$$



Topic : Disk Blocks

Disk block address= 24-bits \Rightarrow no. of blocks = 2^{24}

Size of each block = 2KB

Total Size of disk? = $2^{24} * 2\text{KB} = 2^{35}\text{B} = 32\text{GB}$



Topic : Disk Blocks



Total disk size = 256GB

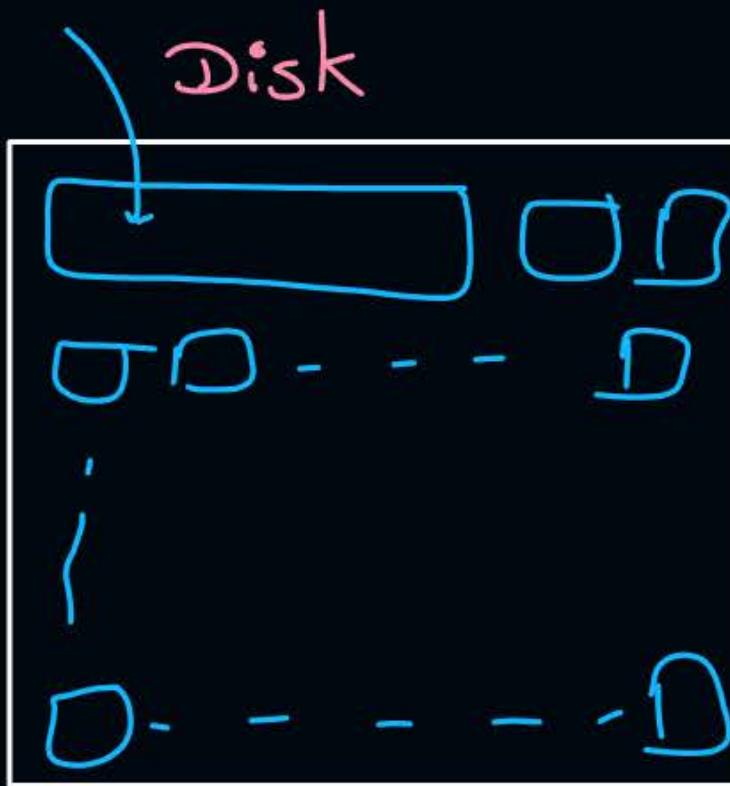
Block Size = 2KB

Disk block address?

$$\text{no. of blocks} = \frac{256\text{GB}}{2\text{KB}} = \frac{2^{38}}{2^{11}} = 2^{27}$$

DBA = 2⁷ bits

some blocks are used to store file directory info



which file stored where and how

ex:-

A disk has 256 blocks each of size 2 kB.
starting 10 blocks are used for storing
directory entries.
Max size of disk remaining to store files ?

soln blocks remaining for file = $256 - 10$
 $= 246$

size for file = $246 * 2 \text{ kB} = 492 \text{ kB}$

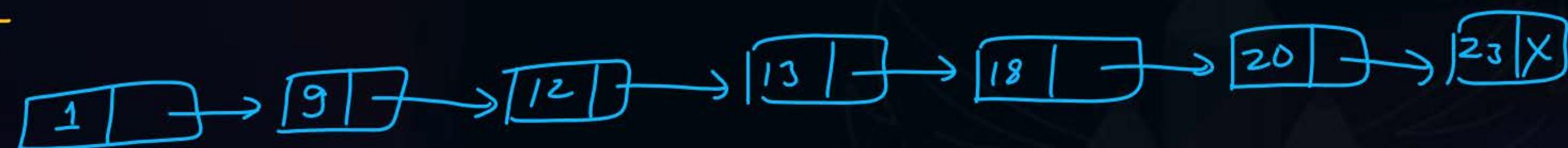


Topic : Free Space Management

1. Free List (linked list of block numbers of all free blocks)
2. Bitmap Method (one bit per block)

e.g:- blocks 1, 9, 12, 13, 18, 20, 23 free

free list:-



Bitmap :- {
 0 => occupied
 1 => free

Block add starts from 0 => 01000000100110000101001



Topic : Free Space Management



1. No searching in free list, but in bitmap we search for first ~~one~~ free block
2. Free list is faster in allocating a free block
3. Free list size is variable, whereas bitmap size is constant



Topic : Question

A particular disk unit uses a bit string to record the occupancy or vacancy of its disk blocks with '0' denoting vacant block and '1' denoting occupied block. A 32-bit part of this string has Hexadecimal value of A3D4B055. The percentage of occupied blocks on the disk for this part is ?

$$\begin{array}{r} 1010\ 0011\ 1101\ 0100\ 10110000\ 0101\ 0101 \\ \hline 2\qquad 2\qquad 3\qquad 1\qquad 3\qquad 0\qquad 2\qquad 2 \end{array}$$

$$\frac{15}{32} * 100 \% = 46.875 \%$$

Ques) A disk has block size = 2 kB

$$\text{no of blocks} = 2^{24}$$

no. of blocks needed to store bitmap for free blocks = 1024 ?

Solⁿ

$$\begin{aligned}\text{no. of bits in bitmap} &= \text{no. of blocks} \\ &= 2^{24} \text{ bits} \\ &= 16 \text{ m bits}\end{aligned}$$

$$\text{no. of blocks} = \frac{16 \text{ m bits}}{2 \text{ kB}}$$

$$= \frac{2^m B}{2 \text{ kB}}$$

$$= 1K$$

$$= 1024$$



Topic : File Allocation Method



describes which disk blocks to be allocated to a file .





Topic : File Allocation Method

1. Contiguous Allocation
2. Linked Allocation
3. Indexed Allocation

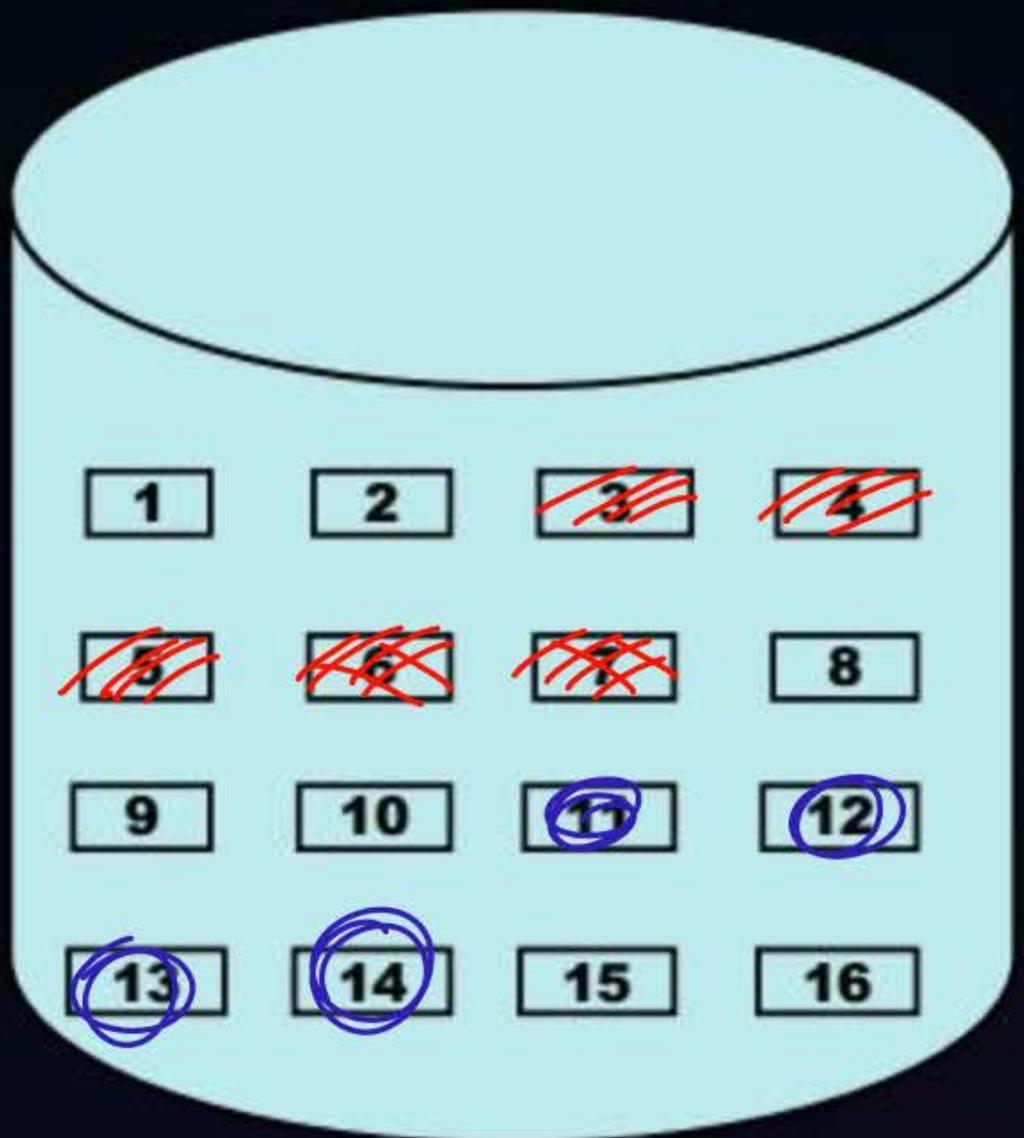




Topic : Contiguous Allocation

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W

Consecutive blocks are used to store file



file allocation

file name	starting block no.	no.of blocks
os.pdf	3	5
COA.pptx	11	4



Topic : Contiguous Allocation



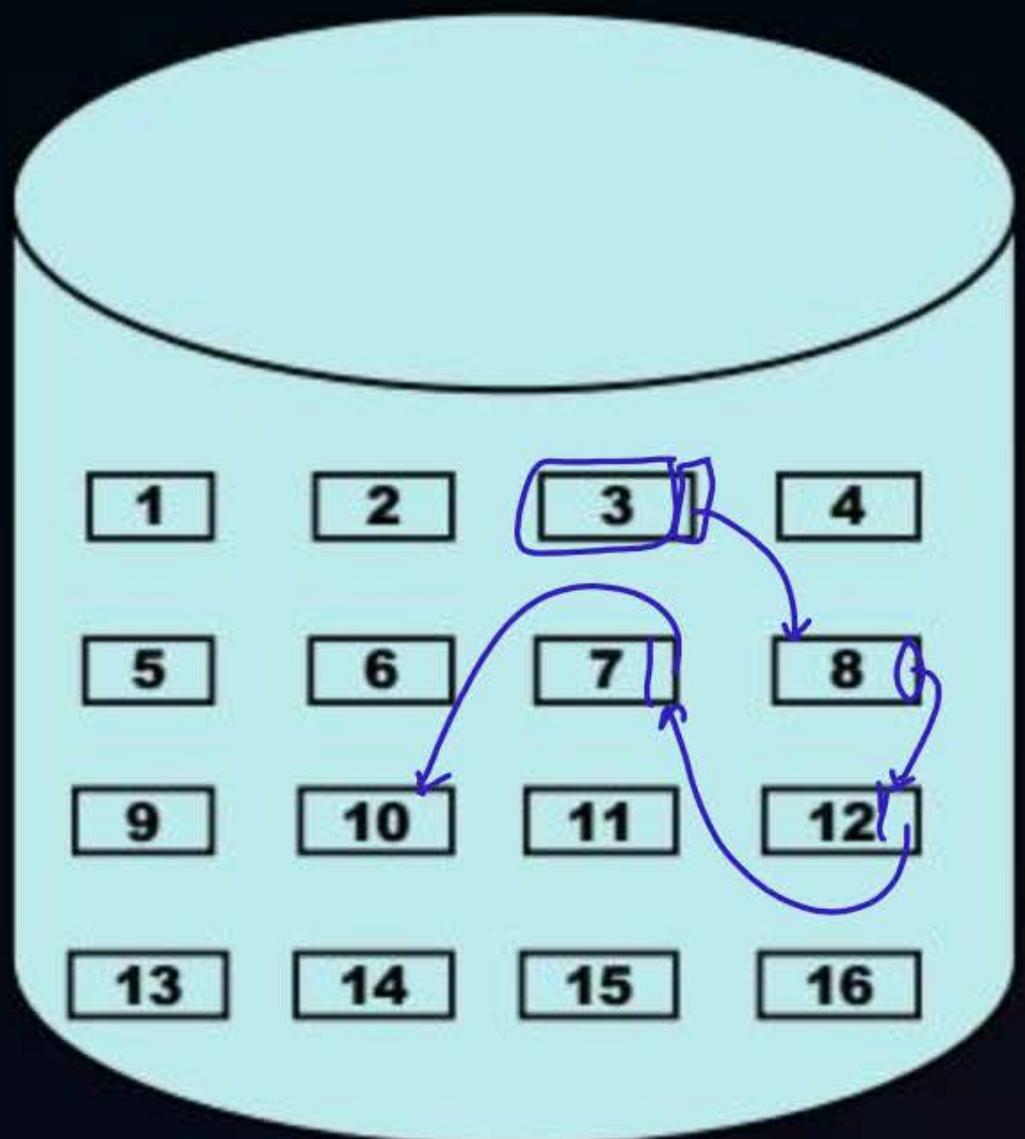
Performance:

1. Fragmentation: Internal, External
2. Increase in File size: Inflexible
3. Type of access: Sequential, Random/direct

Insertion of a new block in file at	no. of block accessed
Beginning	many block accesses needed
Middle	— —
End	few and lesser block accesses



Topic : Linked Allocation



file name	starting block no .	last block no	no. of blocks
coA. pptx	3	10	5



Topic : Linked Allocation

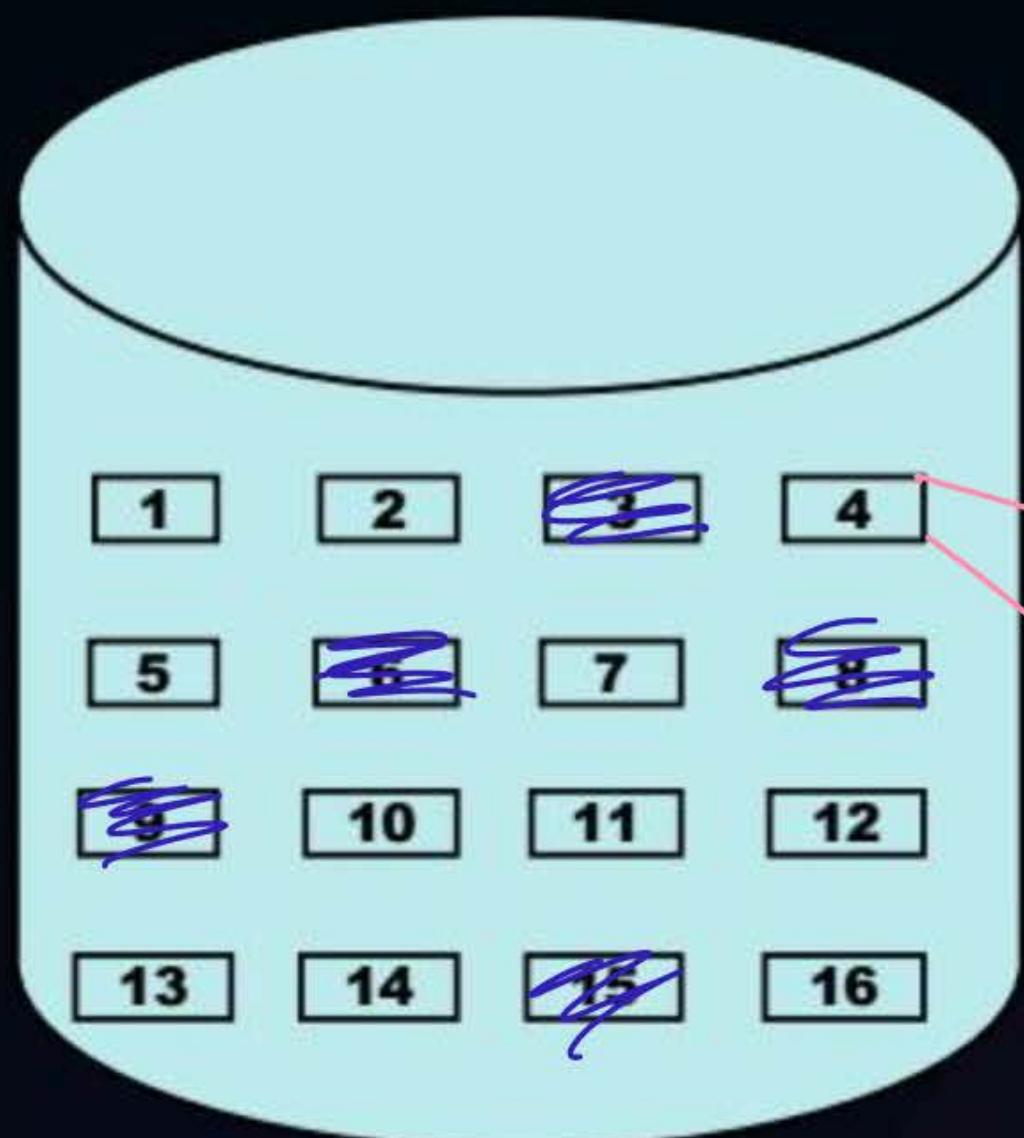
Performance:

1. Fragmentation: Internal
2. Increase in File size: Flexible
3. Type of access: Sequential

Insertion of a new block in file at	no. of blocks accessed
Beginning	very less
Middle	more (as traversal needed)
End	more (— —)



Topic : Indexed Allocation



Allocation table

file name	Index block no.
COA.pptx	4

Index block

6
9
3
8



Topic : Indexed Allocation

Performance:

1. Fragmentation: Internal
2. Increase in File size: Flexible
3. Type of access: Sequential, Random/direct

If file is very big
then
use multi level
indexing

Insertion of a new block in file at	no. of block accesses
Beginning	very less
Middle	- -
End	- -



#Q. Consider two file systems A and B , that use contiguous allocation and linked allocation, respectively. A file of size 100 blocks is already stored in A and also in B. Now, consider inserting a new block in the middle of the file (between 50th and 51st block), whose data is already available in the memory. Assume that there are enough free blocks at the end of the file and that the file control blocks are already in memory. Let the number of disk accesses required to insert a block in the middle of the file in A and B are n_A and n_B respectively, then the value of $n_A + n_B$ is _____?



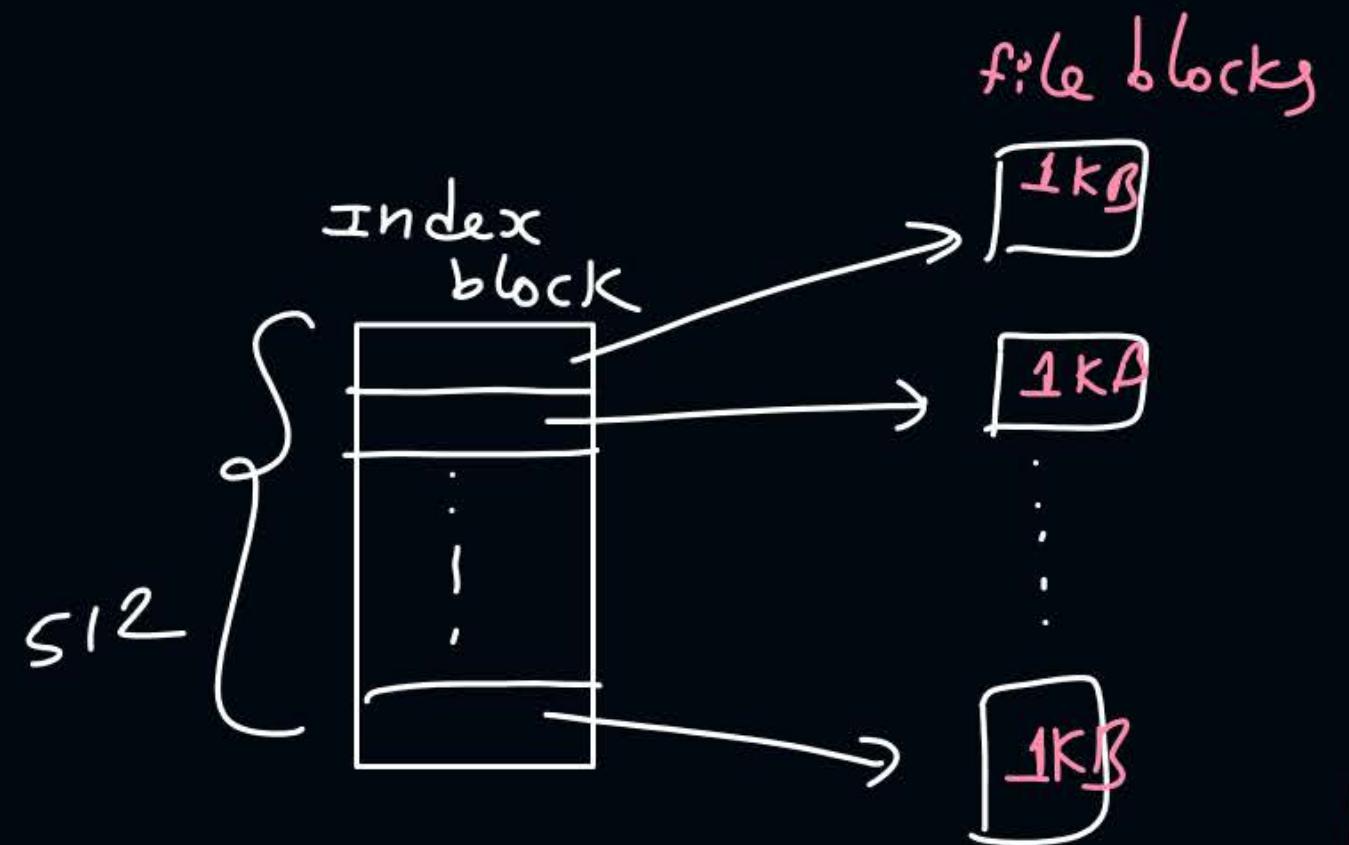
Topic : Question

- Disk block address = 16 bits = 2^B
- Disk block size = 1KB
- Index block = 1KB
- Maximum file size? (single level indexing)

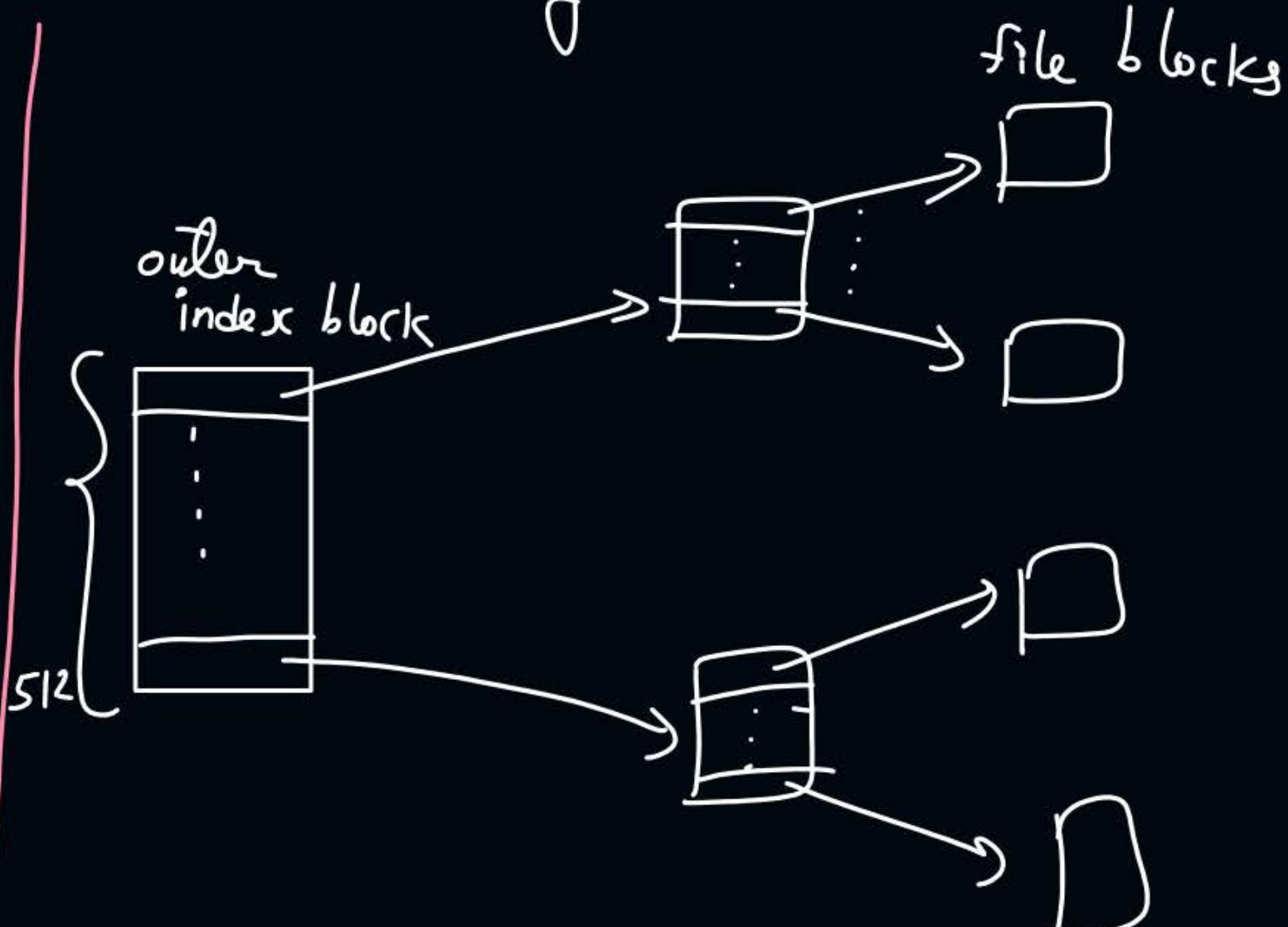
$$\text{no. of disk block numbers per block} = \frac{1KB}{2^B}$$
$$= 512$$

$$\text{max file size} = 512 * 1KB$$
$$= 512 KB$$

1-level indexing



2-level indexing





Topic : Question

- Disk block address = 16 bits = 2^B
- Disk block size = 1KB
- Index block = 1KB
- Maximum file size? (double level indexing)

$$\text{no. of disk block numbers per block} = \frac{1KB}{2B}$$
$$= 512 = 2^9$$

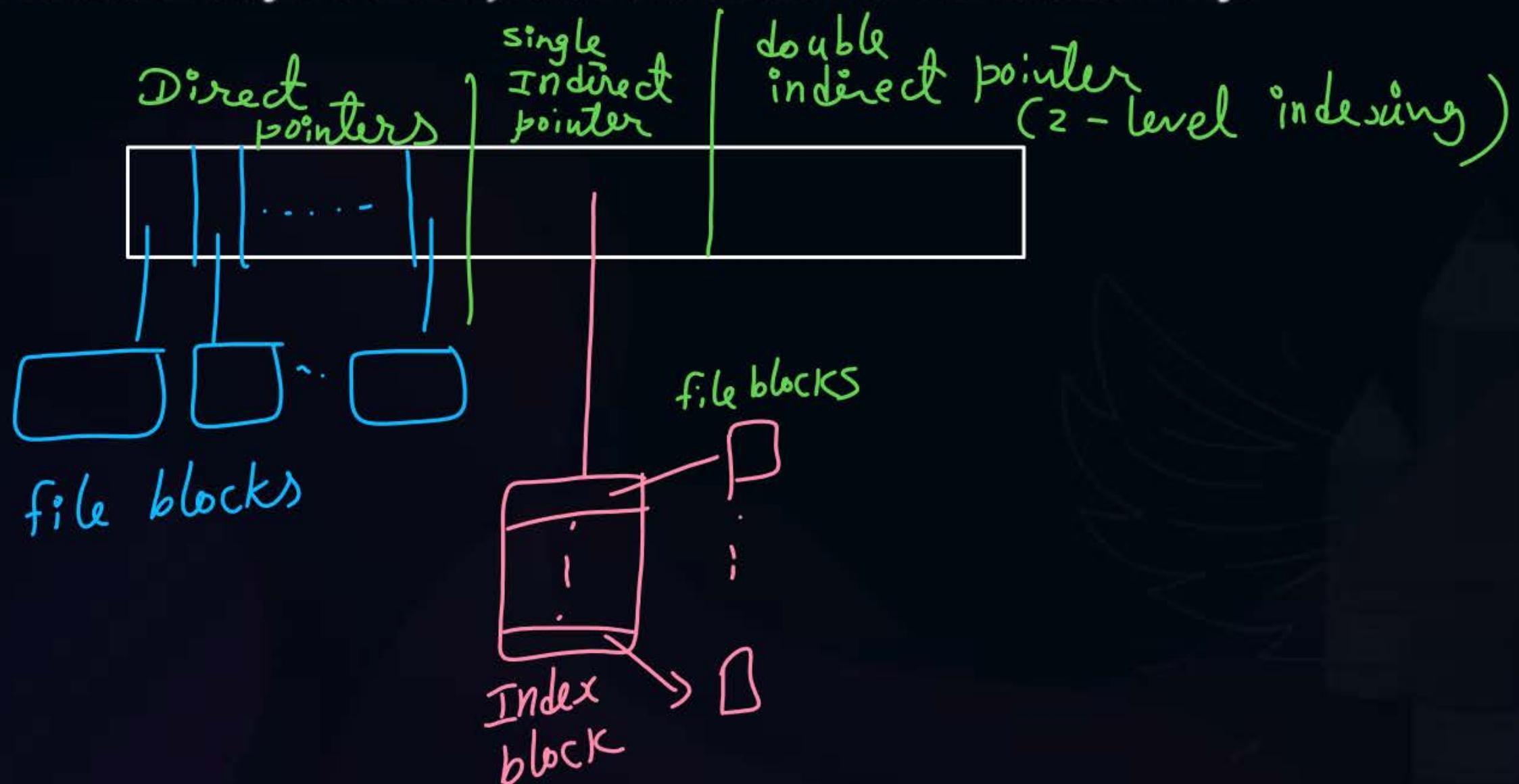
$$\text{max file size} = 2^9 * 2^9 * 1KB = 256MB$$



Topic : Unix I-node Structure



The inode (index node) is a data structure in a Unix-style file system that describes a file-system object such as a file or a directory.





Topic : Unix I-node Structure



The inode (index node) is a data structure in a Unix-style file system that describes a file-system object such as a file or a directory.

- Each inode stores the attributes and disk block locations of the object's data.
- The number of Inode limits the total number of files/directories that can be stored in the file system.



Topic : Question

[GATE-2019]

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W

#Q. The index node (inode) of a Unix-like file system has 12 direct, one single-indirect and one double-indirect pointer. The disk block size is 4 kB, and the disk block addresses 32-bits long. The maximum possible file size is (rounded off to 1 decimal place) 4.0 GB?

4B

$$\text{no. of indexes per block} = \frac{4KB}{4B} = 1K$$

$$\begin{aligned}\text{max file size} &= (12 * 4KB) + (1K * 4KB) + (1K * 1K * 4KB) \\ &= 48KB + 4MB + 4GB \\ &\approx 4.004048GB = 4.04GB\end{aligned}$$



#Q. In a computer system, four files of size 11050 bytes, 4990 bytes, 5170 bytes and 12640 bytes need to be stored. For storing these files on disk, we can use either 100 byte disk blocks or 200 byte disk blocks (but can't mix block sizes). For each block used to store a file, 4 bytes of bookkeeping information also needs to be stored on the disk. Thus, the total space used to store a file is the sum of the space taken to store the file and the space taken to store the book keeping information for the blocks allocated for storing the file. A disk block can store either bookkeeping information for a file or data from a file, but not both.

What is the total space required for storing the files using 100 byte disk blocks and 200 byte disk blocks respectively?

A

35400 and 35800 bytes

B

35800 and 35400 bytes

C

35600 and 35400 bytes

D

35400 and 35600 bytes



Topic : Solution

File Size	No. of blocks to store file	No. of blocks to store bookkeeping	Total blocks	Total Size
11050 Bytes				
4990 Bytes				
5170 Bytes				
12640 Bytes				



Topic : Solution

File Size	No. of blocks to store file	No. of blocks to store bookkeeping	Total blocks	Total Size
11050 Bytes				
4990 Bytes				
5170 Bytes				
12640 Bytes				



2 mins Summary



Topic File System

Topic Blocks and File Allocation Methods

Topic Disk Scheduling

Topic SSTF, Scan, Look



Happy Learning

THANK - YOU