



# CS & DA



## Database Management System

### Transaction and Concurrency Control

DPP 01 (Discussion Notes)



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#Q. How many serial schedules can be formed with 4 transactions.

1)  $\tau_1 \quad \tau_2 \quad \tau_3 \quad \tau_4$

2)  $\tau_1 \quad \tau_2 \quad \tau_4 \quad \tau_3$

3)  $\tau_1 \quad \tau_4 \quad \tau_2 \quad \tau_3$

$$4! = \underline{24}$$

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.

.

#Q. How many concurrent schedules can be formed with 3 transactions having 4, 3 & 2 operations respectively.

$$\tau_1 = 4 \text{ ops}$$

$$\tau_2 = 3 \text{ ops}$$

$$\tau_3 = 2 \text{ ops}$$

$$\overline{\text{Total ops} = 9 \text{ ops}}$$

$$\frac{9!}{4! 3! 2!} = 1260$$

#Q. Consider the following schedule

S:  $R_1(A); R_3(A); R_2(A); W_1(B); R_2(B); R_3(A); W_2(C); R_3(C)$  over the transaction  $T_1, T_2 \& T_3$ .

If transaction  $T_1$  fails just after  $R_3(C)$  by transaction  $T_3$ , then which transaction need to be rolled back along with  $T_1$ .

A  $T_2$

B  $T_3$

$T_1$   $T_K$   
 $\omega(x)$

C Both  $T_2 \& T_3$

D None

$R(x)$

$T_2$

#Q. Consider the following transactions.

$T_1; W_1(A); W_1(B); R_1(C); C_1;$

$T_2; W_2(B); R_2(B); C_2;$

How many schedules of  $\underline{T_1}$  &  $\underline{T_2}$  are irrecoverable.

$$\checkmark \frac{T_1 \quad T_2}{\omega(x)}$$

$$\frac{T_1 \quad T_2}{\omega(x)}$$

$$\frac{R(x)}{C_2}$$

$C_2$

$C_1$

$$\left\{ \begin{array}{c} \frac{T_1 \quad T_2}{\omega(B)} \\ \omega(A) \\ \omega_1(B) \end{array} \right. - 2 \text{ ways}$$

$$\left. \begin{array}{c} \frac{R(B)}{C_2} \\ C_2 \\ C_1 \end{array} \right\} 3 \text{ ways}$$

$$2 \times 3 - \underline{6 \text{ ways}}$$

#Q. Two schedules  $S_1$  and  $S_2$  are called conflict equivalent if  $S_1$  can be derived from  $S_2$  by a sequence of swaps of non-conflicting operations.

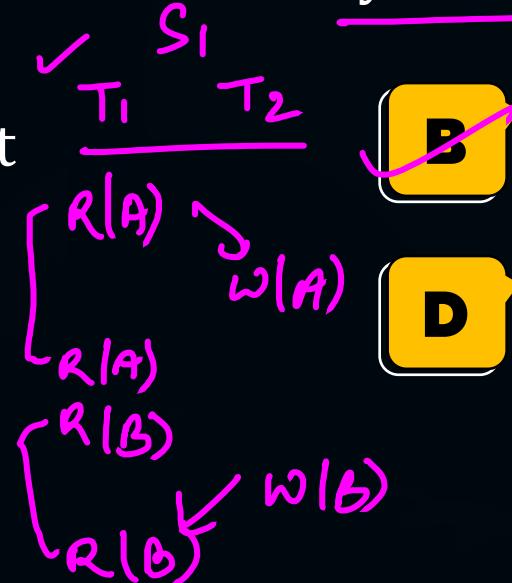
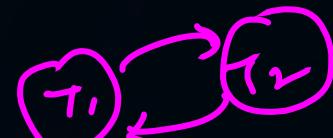
Consider the two statements:

I → If two schedule are conflict equivalent , then their precedence graphs are identical.

✗ II → If two schedules involve same set of transactions, and their precedence graphs are identical. Then they are conflict equivalent, ✗

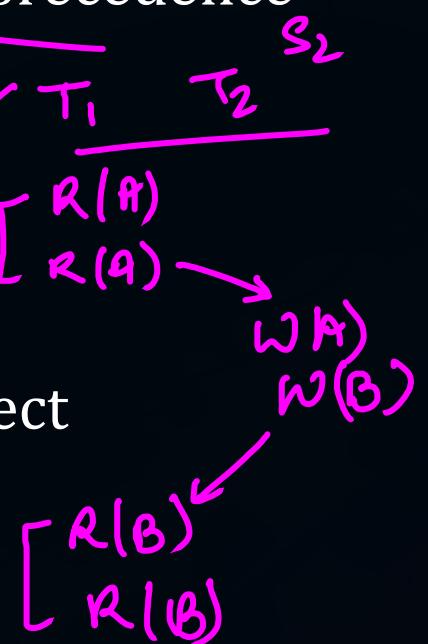
**A** Both I & II are correct

**C** Only II is correct



Only I is correct

Neither I nor II is correct

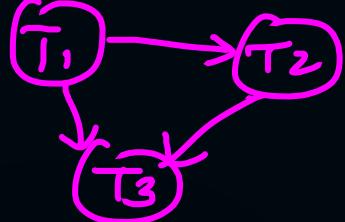


#Q. Which of the following schedules is/are irrecoverable.

- X **A** R<sub>1</sub>(A), R<sub>2</sub>(C), R<sub>1</sub>(C), R<sub>3</sub>(A), R<sub>3</sub>(B), W<sub>1</sub>(A), C<sub>1</sub>, W<sub>3</sub>(B), C<sub>3</sub>, R<sub>2</sub>(B), W<sub>2</sub>(C), W<sub>2</sub>(B), C<sub>2</sub>
- B** R<sub>1</sub>(A), R<sub>2</sub>(C), R<sub>1</sub>(C), R<sub>3</sub>(A), R<sub>3</sub>(B), W<sub>1</sub>(A), W<sub>3</sub>(B), R<sub>2</sub>(B), W<sub>2</sub>(C), W<sub>2</sub>(B), C<sub>1</sub>, C<sub>2</sub>, C<sub>3</sub>
- X **C** R<sub>1</sub>(A), R<sub>2</sub>(C), R<sub>3</sub>(A), R<sub>1</sub>(C), R<sub>2</sub>(B), R<sub>3</sub>(B), W<sub>1</sub>(A), C<sub>1</sub>, W<sub>2</sub>(C), W<sub>3</sub>(B), W<sub>2</sub>(B), C<sub>3</sub>, C<sub>2</sub>
- X **D** All are recoverable

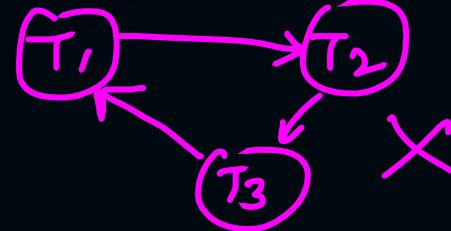
#Q. Which of the following schedules is/are conflict serializable.

**A**



R<sub>1</sub>(x), W<sub>1</sub>(y), R<sub>2</sub>(y), W<sub>2</sub>(z), R<sub>3</sub>(z), W<sub>3</sub>(x)

**B**



W<sub>3</sub>(x), R<sub>1</sub>(x), W<sub>1</sub>(y), R<sub>2</sub>(y), W<sub>2</sub>(z), R<sub>3</sub>(z)

**C**



R<sub>1</sub>(x), R<sub>2</sub>(x), W<sub>1</sub>(y), W<sub>2</sub>(y), R<sub>1</sub>(y), R<sub>2</sub>(y), W<sub>2</sub>(z)

**D**

R<sub>1</sub>(x), R<sub>2</sub>(x), R<sub>1</sub>(y), R<sub>2</sub>(y), R<sub>3</sub>(x), W<sub>1</sub>(x), W<sub>2</sub>(y)

**D**



T<sub>3</sub>

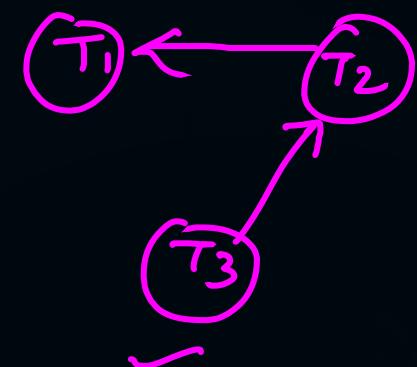
X

#Q. Consider the following schedule S.

Schedule S is conflict equivalent to which of the following serial schedule.

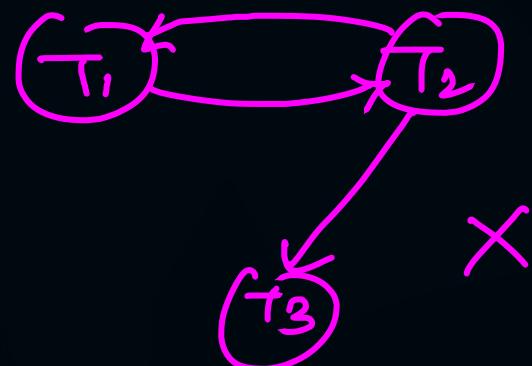
- A**  $T_1 \rightarrow T_3 \rightarrow T_2$
- B**  $T_3 \rightarrow T_1 \rightarrow T_2$
- C**  $\cancel{T_3 \rightarrow T_2 \rightarrow T_1}$
- D**  $T_2 \rightarrow T_1 \rightarrow T_3$

| S        |          |          |
|----------|----------|----------|
| $T_1$    | $T_2$    | $T_3$    |
| $R_1(x)$ |          |          |
| $W_1(x)$ |          |          |
|          | $R_2(x)$ |          |
|          | $R_2(z)$ |          |
|          | $R_2(y)$ |          |
|          | $W_2(y)$ |          |
| $W_1(z)$ |          | $R_3(y)$ |



#Q. Consider the following schedule S.

| S                  |                    |                    |
|--------------------|--------------------|--------------------|
| T <sub>1</sub>     | T <sub>2</sub>     | T <sub>3</sub>     |
|                    | R <sub>2</sub> (B) |                    |
| R <sub>1</sub> (A) | W <sub>2</sub> (A) |                    |
| W <sub>1</sub> (B) |                    | R <sub>3</sub> (A) |
|                    | W <sub>2</sub> (B) |                    |
|                    |                    | W <sub>3</sub> (B) |



Initial read

T<sub>2</sub> ✓

Updated read

T<sub>2</sub> → T<sub>1</sub>

T<sub>2</sub> → T<sub>3</sub>

Which of the following options is/are correct ?

X **A**

The schedule is conflict serializable schedule

✓ **B**

The schedule is view serializable schedule

X **C**

$T_2 \rightarrow T_1 \rightarrow T_3$  is conflict equivalent serial schedule to s.

✓ **D**

$T_2 \rightarrow T_1 \rightarrow T_3$  is view equivalent serial schedule to s.

#Q. Consider the following schedule S.

S:  $R_1(A)$ ,  $W_2(B)$ ,  $R_2(C)$ ,  $W_3(B)$ ,  $W_2(A)$ ,  $W_1(A)$ ,  $R_3(B)$ ,  $R_1(A)$ ,  $R_2(C)$ ,  $R_3(C)$ ,  
 $W_2(C)$ ,  $C_1$ ,  $C_3$ ,  $C_2$ ,

Schedule S suffers from which of the following problems.

| $T_1$  | $T_2$ | $T_3$ |
|--------|-------|-------|
| $R(A)$ |       |       |

$w(B)$   
 $R(c)$

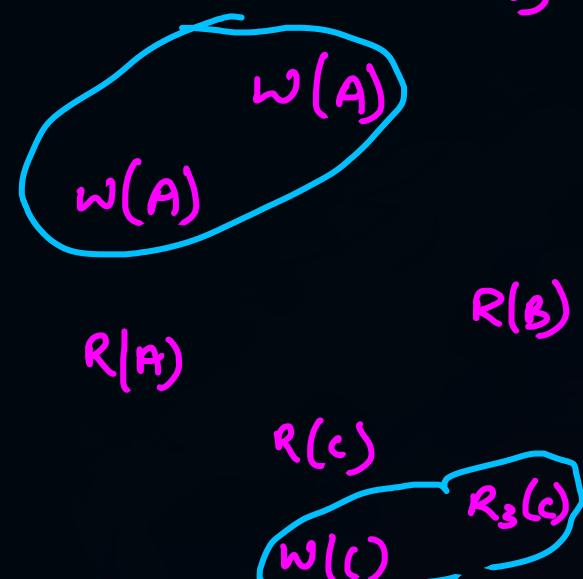
$\downarrow$   
 $w(B)$

**A** Irrecoverability

**B** Cascading Roll back

**C** Lost update problem

**D** RW Problem





c<sub>1</sub>

c<sub>3</sub>

c<sub>2</sub>

THANK - YOU