CS302: Paradigms of Programming

Lab 2: Church Numerals

March 9th, 2021

Recall the claim the instructor had made that everything apart from lambdas is syntactic sugar. Let us get a flavour of the same.

Say we define the first number zero as:

```
(define zero (lambda () '()))
```

That is, zero is represented just as a lambda that returns an empty list.

Say we additionally have the following two definitions for computing the successor and the predecessor of a given number x:

```
(define (succ x) (lambda () x)) ; i.e., wrap x in a lambda
(define (pred x) (x)) ; i.e., unwrap x by calling that lambda
```

The purpose of the above two clever definitions is to ensure the property that the predecessor of the successor of a number n is n. Make sure you see this in the definitions (in terms of when is a procedure applied versus defined, etc.).

What is the number one? Successor of zero! What is the number two? Successor of one (or successor of successor of zero)! And so on. Let's define a few numbers this way:

```
(define one (succ zero))
(define two (succ one))
(define three (succ two))
(define four (succ three))
(define five (succ four))
```

Now say we have a way to compute if something is zero:

```
(define (is-zero? x) (null? (x)))
(works because null? checks for empty list.)
```

Next, let us define a procedure that checks if two of our newly defined *Church numerals* are equal:

```
(define (is-equal? x y)
  (cond ((is-zero? x) (is-zero? y))
        ((is-zero? y) (is-zero? x))
        (else (is-equal? (pred x) (pred y)))))
```

Basically, as the only actual equality check we have is for zeros (in the form of is-zero?), what we are doing in (is-equal? x y) is getting down to using is-zero? by recursively reducing x and y using the pred function.

Write code to verify the following kinds of properties:

- 1. zero is not equal to one, but is equal to zero
- 2. four is equal to succ(succ(succ(succ(zero))))
- 3. The predecessor of the successor of two is two

Apart from checking for equality, what else are numbers useful for? Addition, subtraction, multiplication, and so on. Understand and programmatically verify (like the items above) that the following add-church procedure works:

```
(define (add-church x y)
  (if (is-zero? y)
    x
    (add-church (succ x) (pred y))))
```

Now write your own versions of subtract-church and multiply-church.

Absorb the fact that we were able to implement *numbers* and (some of) the important associated operations just with *lambdas*; that should give you an idea of the power of the **lambda calculus** (and its sugarless world:p).

If you found this interesting, go ahead and look up how could you define booleans (a smaller set than numbers, isn't it?) and operations over booleans using lambdas, leading to the notion of *Church booleans*.