

The Power of Two-Choices in reducing discrepancy

Ohad N. Feldheim

Based on joint works with:

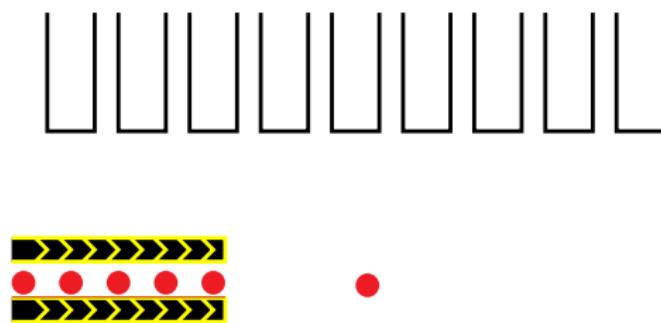
Ori Gurel-Gurevich (HUJI)

Raaz Dwivedi & Aaditya Ramdas (Berkeley)

June 2017

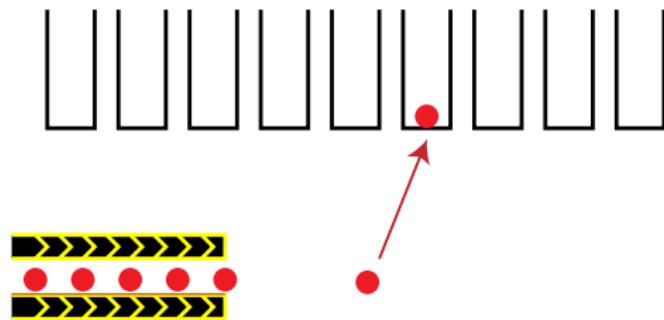
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Consider an online process in which N balls are randomly assigned, one by one to M bins.



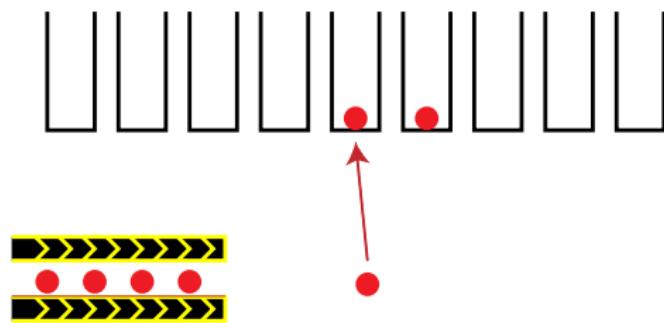
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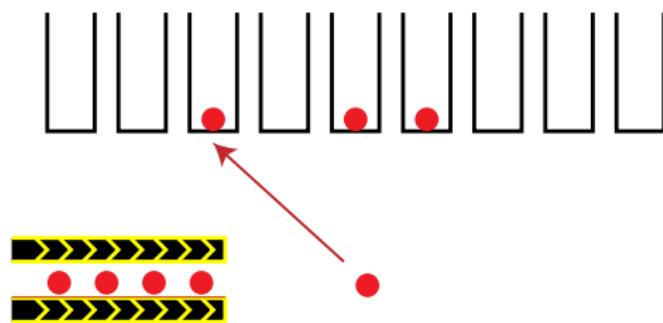
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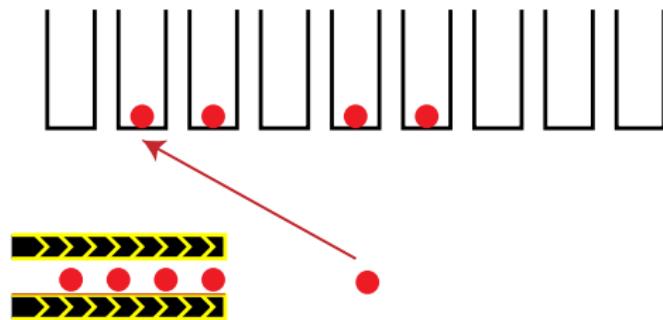
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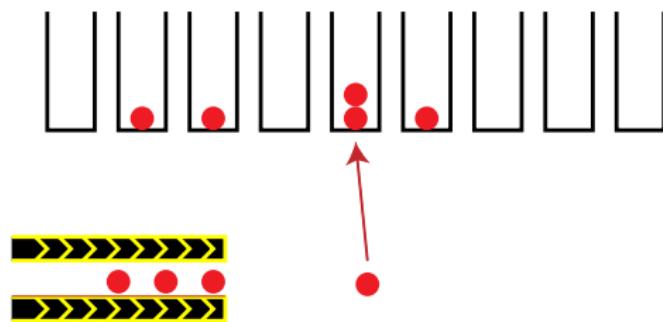
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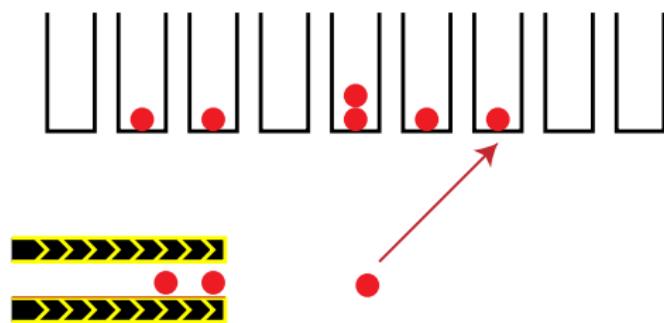
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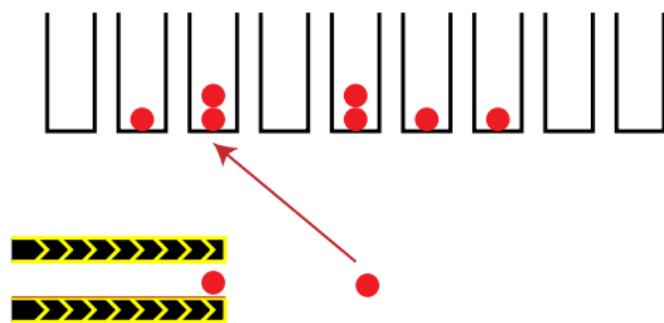
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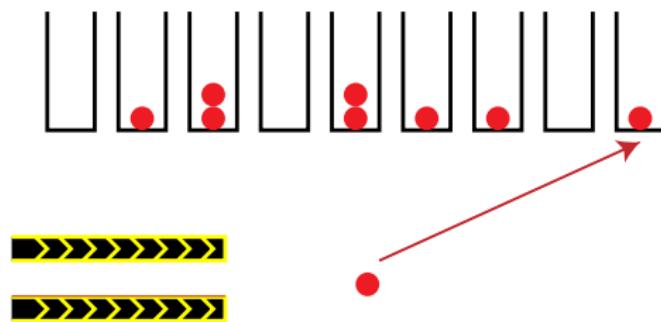
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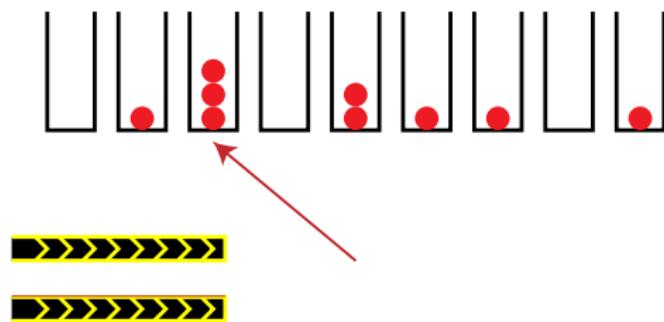
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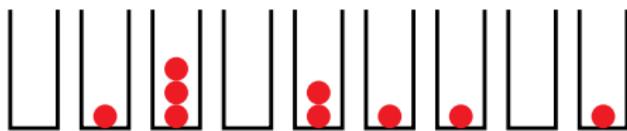
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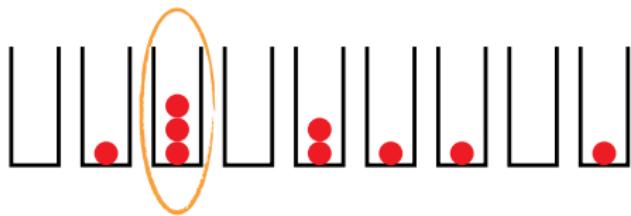
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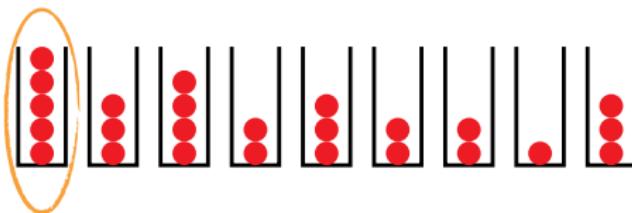
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typical deviation from expectation is $\Theta\left(\sqrt{\frac{N}{M}}\right)$
- *Load balancing* is an effort to reduce these quantities.
(possible with control over the distribution of the balls.)

Power of two choices

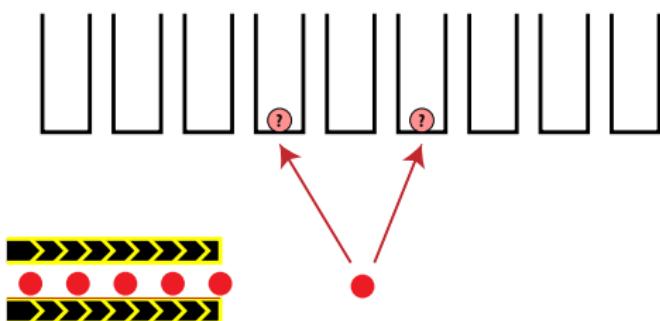
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Greedy strategy: choose the least occupied cell.

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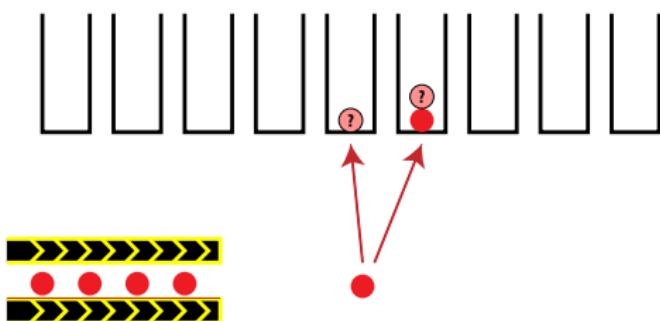
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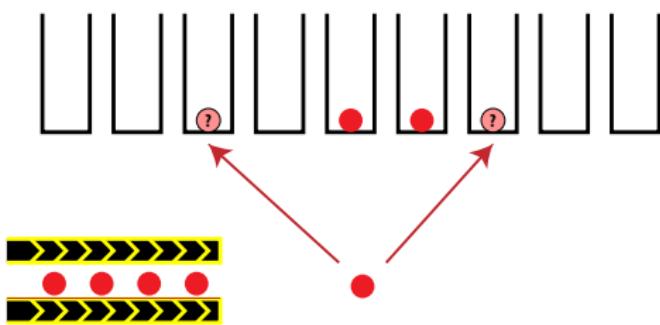
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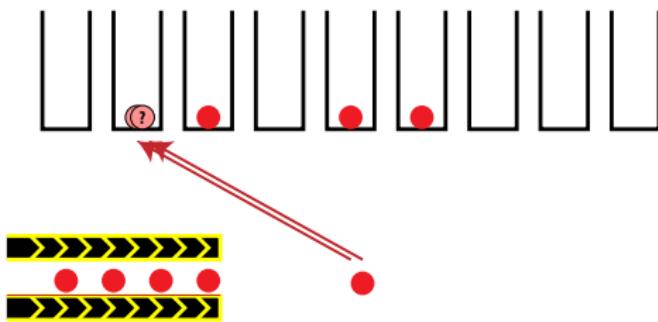
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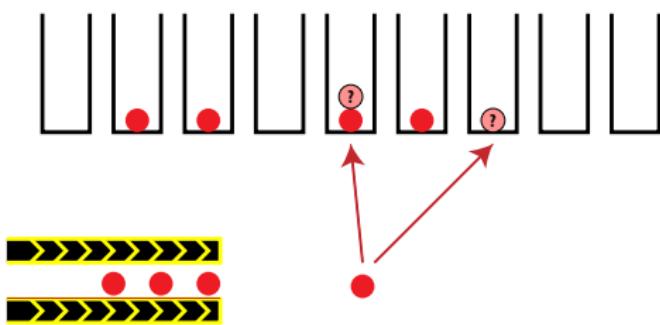
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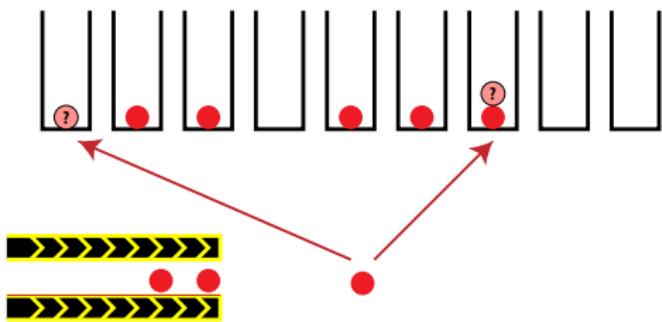
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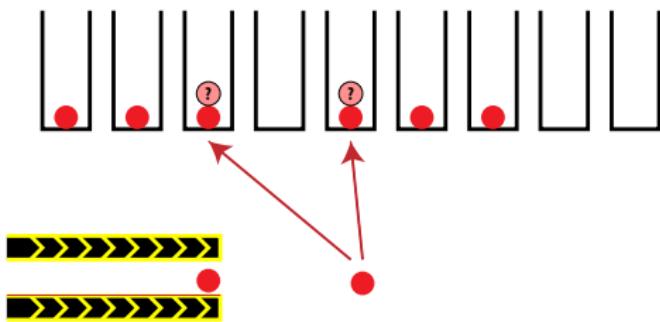
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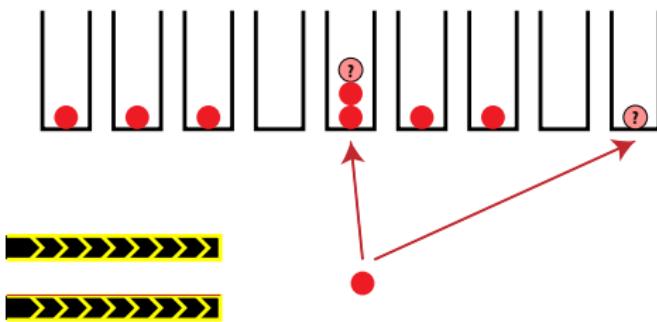
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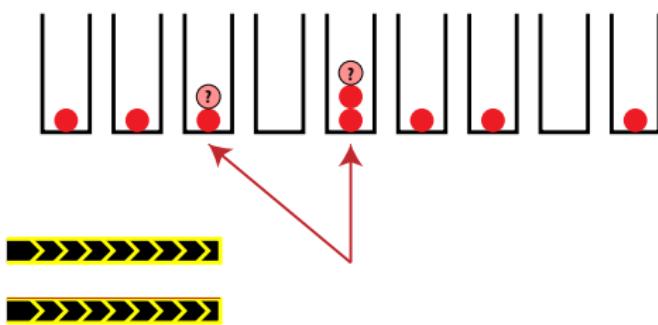
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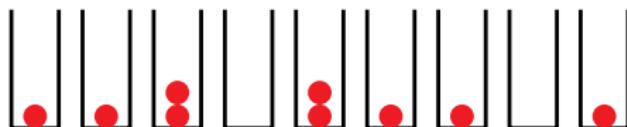
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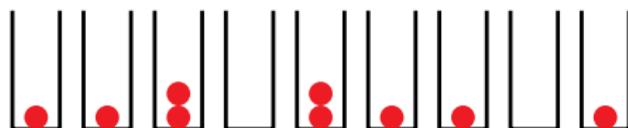
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Balls	no-choice max	2-choices max	no-choice typ. dev.	2-choices typ. dev.
M	$\frac{\log M}{\log \log M}$	$\frac{\log \log M}{2}$	$O(1)$	$O(1)$
$N \gg M$	$\Theta\left(\sqrt{\frac{N \log M}{M}}\right)$	$\log M$	$\Theta\left(\sqrt{\frac{N}{M}}\right)$	$O(1)$

Power of two choices - remarks

This observation had many applications

- Server load-balancing
- Distributed shared memory
- Efficient on-line hashing
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It is interesting to note that

- More choice does not yield significant improvement.
- If balls keep appearing and dying at rate 1 the phenomenon persists (Luczak & McDiarmid '05)
- However if one can't keep track of the number of balls per bin (due to having $M^{1-\epsilon}$ bits of memory), then no asymptotic improvement over no-choice is possible (Alon, Gurel-Gurevich, Lubetzky '09)

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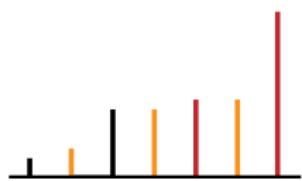
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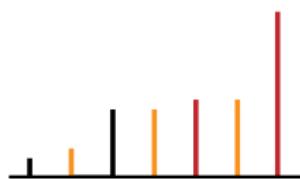
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Two-choices and interval partitions

Benjamini: Can two choices regulate interval partitions?



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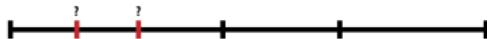


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The empirical interval distribution:



Uniform interval partition $\rightarrow \text{Exp}(1)$
Kakutani interval partition $\rightarrow \mathcal{U}(0, 2)$
Max-2 interval partition $\rightarrow ???$

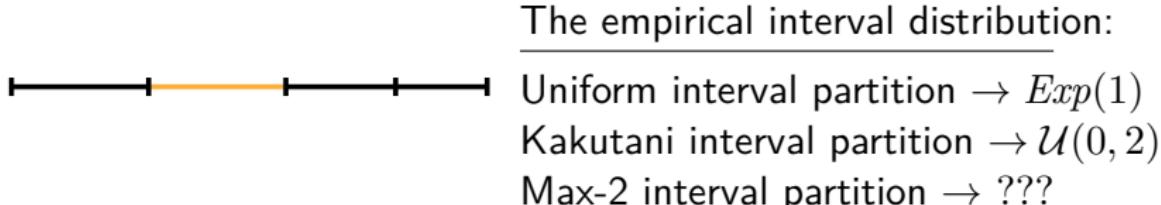
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However, even Kakutani offers no significant improvement in terms of interval variation.

Convergence of 2-Max interval partition process

Studying 2-Max is a rather difficult task:

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In a sense, corresponds to Alon, Gurel-Gurevich, Lubetzky.



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We show that a global strategy can reduce discrepancy significantly.

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Power of two-choices in regulating interval discrepancy (F. & Gurel-Gurevich 2016+)

In a power of two-choices process on $\mathcal{U}[0, 1]$, the chooser can obtain

$$\lim_{n \rightarrow \infty} \Pr \left(\max_{a,b} |\mu^n((a, b)) - \mu((a, b))| < \frac{C \log^3 N}{N} \right) = 1.$$

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- Cf. lower bound of $\frac{C \log N}{N}$, no choice estimate $\sqrt{\frac{C \log N}{N}}$.

Global strategy that regulates interval variation

Discrete counterpart

For N balls on $\mathcal{U}([M])$, a probabilistic retry strategy obtains

$$\Pr \left(\max_{a < b \in [M]} |\mu^n([a, b]) - \mu([a, b])| > \Delta \log^3 M \right) \leq Ce^{-c\Delta}.$$

Stochastic point of view on the power of two choices



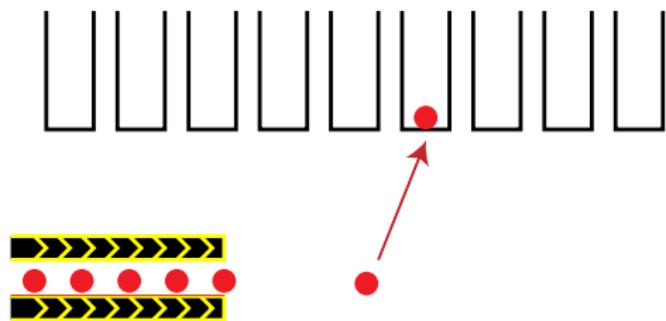
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One-retry: A related intermediate setup. The chooser is only offered a chance to re-roll the target bin once per ball.



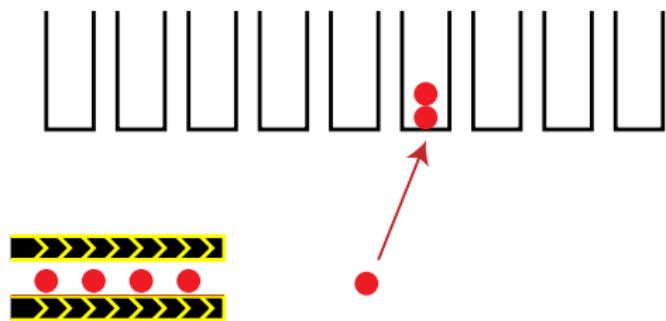
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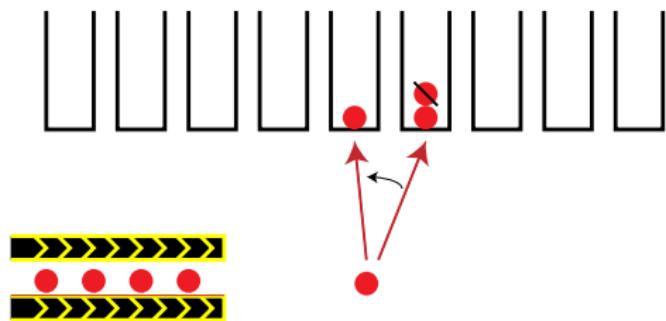
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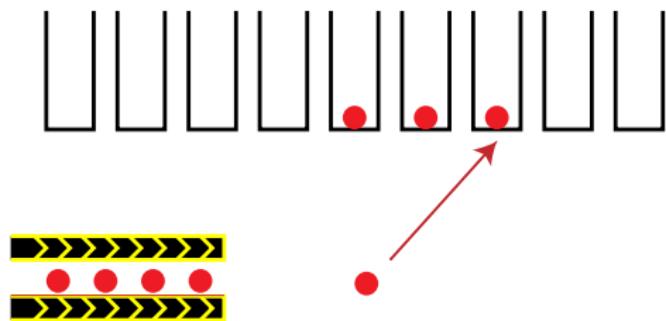
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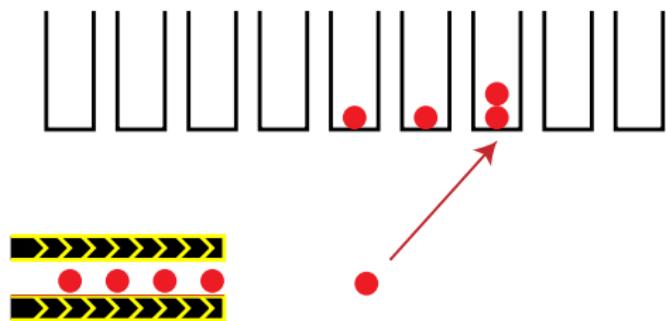
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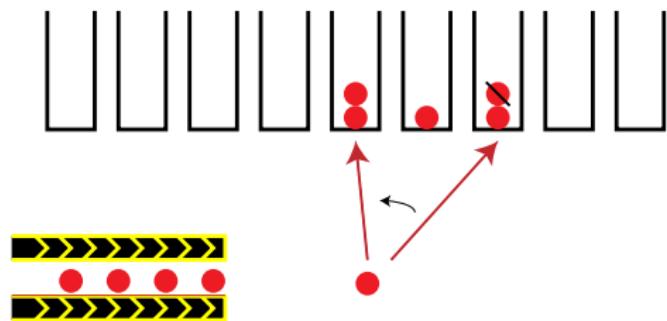
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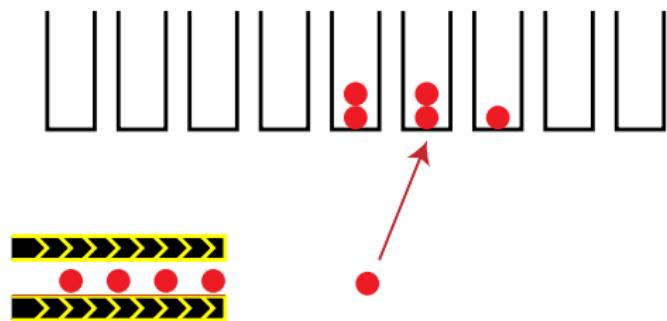
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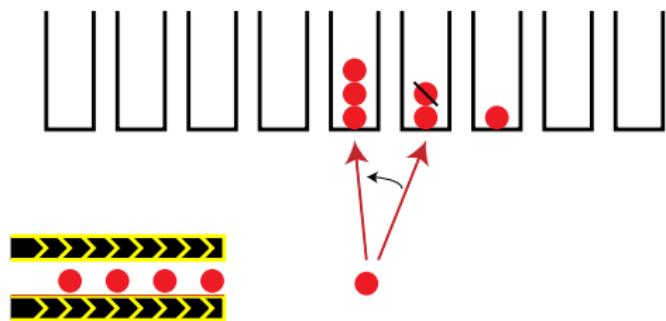
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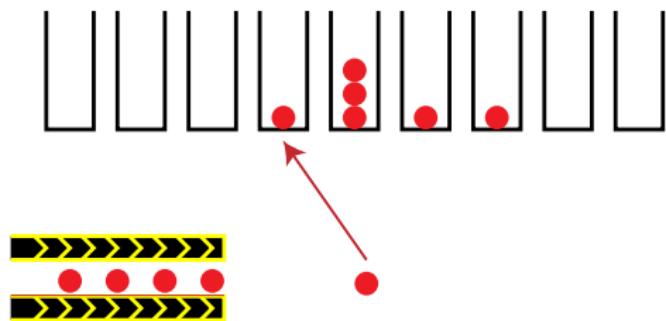
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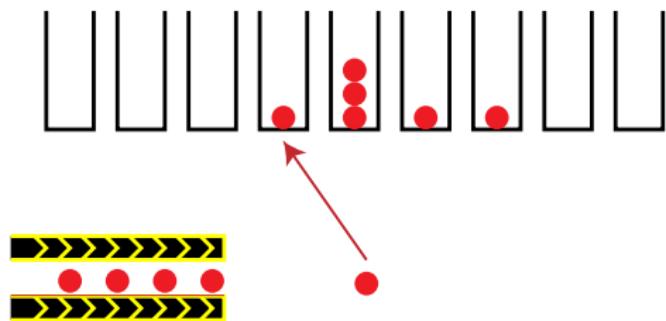
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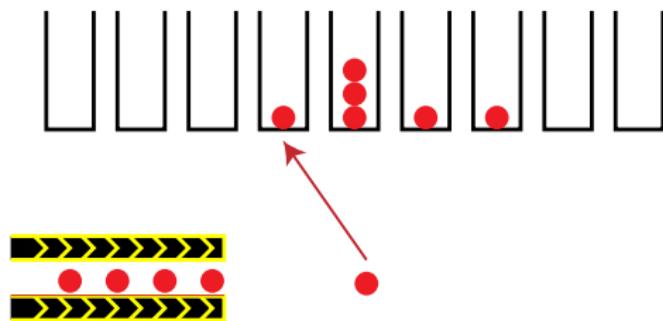
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STOCHASTIC APPROACH TO TWO-CHOICES: REGULATING BALLS AND BINS

Large family of one-retry distributions

What kind of distributions could be realized using one retry?

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Every distribution whose density $\frac{1}{2M} < g(x) < \frac{3}{2M}$ on $[M]$ could be realized by a (probabilistic) one-retry strategy.

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The probability of a random bin to be re-rolled is $\sum_{i=1}^M \frac{f(x)}{M} = \frac{1}{2}$. Hence the probability that x is chosen is now

$$\frac{1-f(x)}{M} + \frac{1}{2M} = g(x)$$

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i.e. such a process has typical fluctuation $O(\theta)$.

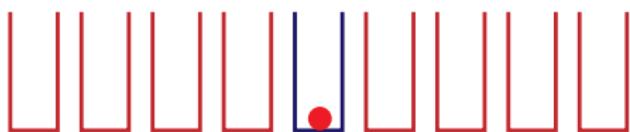
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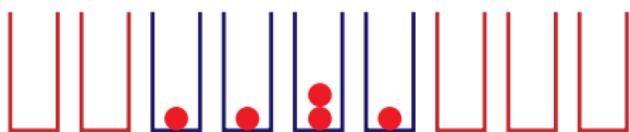
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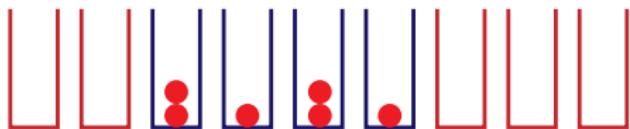
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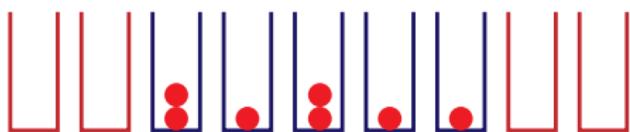
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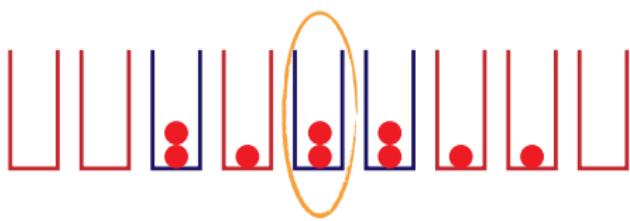
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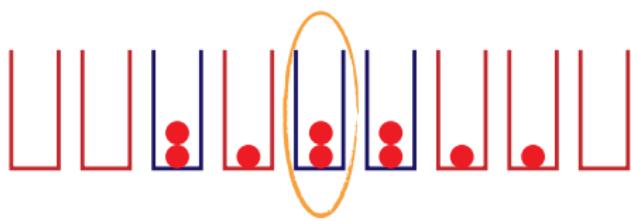
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Among M such processes, the extremal fluctuation is $O(\theta \log M)$.

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All that is left is to show that the same holds after N balls were distributed. - i.e. show concentration of the stopping time.



STOCHASTIC APPROACH TO TWO-CHOICES: REGULATING INTERVAL PARTITION



Balancing poisson point process

Consider two point processes X_t^0, X_t^1 with changing causal intensities $\lambda^0(t), \lambda^1(t)$ (which may depend on other variables), which satisfy

$$\lambda_t^1, \lambda_t^0 < 2 \quad \forall t$$

$$\lambda_t^0 - \lambda_t^1 \geq \theta \quad X_t^0 \leq X_t^1$$

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We now define an intensity μ_t on $[0, 1]$ which could be realized by a one-retry strategy. We will show that this intensity regulates all diadic intervals.

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Each $(a, b) \subset [0, 1]$ has $(a, b) = \sum_{i=0}^{\lfloor \log_2 N \rfloor} I_i + R$ where I_i are diadic and $|R| < 1/N$.

Hierarchy of drifts

We build a hierarchy of $\log_2 N$ drifts of strength $C/\log N$, to control diadic discrepancies.

Hierarchy of drifts

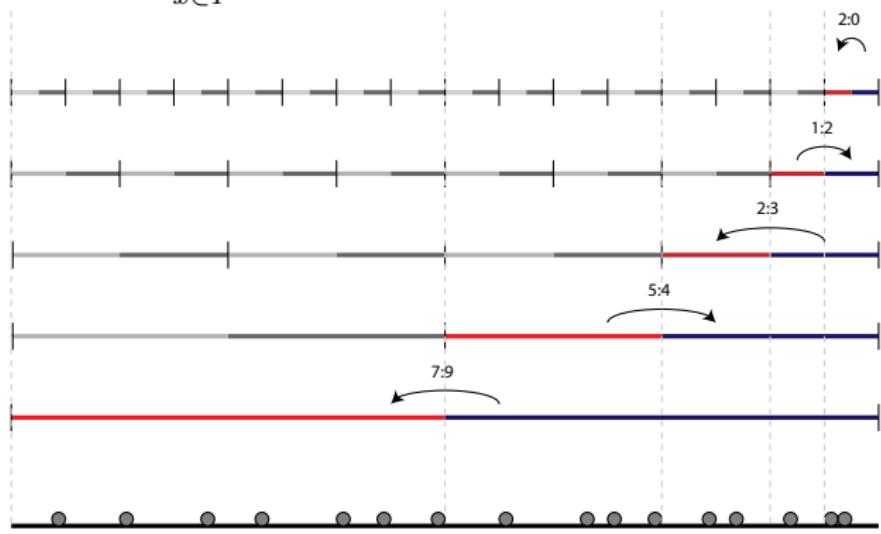
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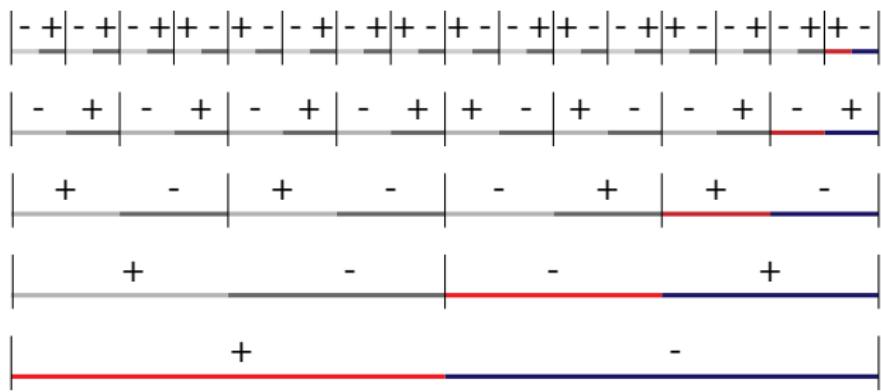
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Two-choices for Balls and Bins
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Interval Partition
○○○

Results and Method
○○○○

Method - Balls & Bins
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Method - Interval Partition
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Future directions

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- Other spaces

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- Other measures of discrepancy.

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- Reducing the power of the log.

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THANK YOU.