ECS 158 Final Project: An Attempt At Parallelizing R's Phylobase::ShortestPath()

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Abstract

This report attempts paralleize CRAN's phylobase package's shortestPath function in RSnow, OpenMP and CUDA. The function takes in a phylogeneic tree, two nodes in the tree and produces the shortest path of nodes inbetween them. The RSnow implementation built on top fo the existant code and paralleized the descendants function. The OpenMP implementation did a similar approach, except with; ALICIA WHAT DID YOU DO;. The CUDA implementation took the brute force approach to the shortestPath problem. Overall,; SENTENCE ABOUT TEST RESULTS;; CLOSE WITH A MORE GENERAL STATEMENT BUILDING OFF RESULTS OF THE TESTS;

1 Introduction and Motivation

Alongside the rise of "big data" in the recent years, bioinformatices has gained considerable momentum. But the a consistent issues remain: what do we do with all the data and how do we make sense of it at a reasonable rate? The R community has taken a stab at those issues. For this report, we are examining the R's phylobase package, which provides the base class and functions for phylogenetic or evolutionary structures and comparitive data. (CITE) We will be focusing our efforts in making the "treewalk" utility functions, such as finding descendants/ancestors and shortest pathes, fast through three different parallel programming models (RSnow, OpenMP and CUDA).

2 Approaches

Only pseudocode or key chunks of code of each implementation are shown or described, see Appendix A for code details.

2.1 Original

The original R implementation calculated the shortest path between the two nodes of interest by first calculating their Most Recent Common Ancestor (MRCA). Then that MRCA's descendants are calculated and compared to the two nodes of interest's ancestors. Any overlap is stored and that is the shortest path.

The C version of the descendants function, called Cdescendant(), works by first marking given node in a preordered list of edges. Then the direct descendants (or children) of the given node is marked. Cdescendant() then iterates through the other edges and marks each marked node's direct descendants (children).

The C version of the ancestors function works the same as Cdescendants, expect direct ancestors (or parents) are marked instead of direct descendants (children).

Pseudocode (source code in original phylobase package):

```
descendants (tree, given_node) {
 2
      #let x be the edges of tree listed in PREORDER,
 3
      #with the older node occupying the first column
      # C function call
 4
 5
      is Descendant \leftarrow Cdescendants (x[,1], x[,2], given_node)
 6
      retval <- getNode(tree, isDescendant)
 7
 8
 9
   ancestors (tree, node1) {
      #let x be the edges of tree listed in POSTORDER,
10
      #with the older node occupying the first column
11
12
      # C function call
      isAncestor \leftarrow Cancestors(x[,1], x[,2], given_node)
13
14
      retval <- getNode(tree, isAncestor)</pre>
15
   }
16
   MRCA(tree, node1, node2 ... noden) {
17
18
       nodes <- unique (node1, node2, ..., noden)
19
       ancests <- lapply (nodes, ancestors, phy=phy, type="ALL")
20
        retval <- getNode(phy, max(Reduce(intersect, ancests)))
21
22
   shortestPath(tree, node1, node2){
23
         t1 <- getNode(tree, node1)
24
         t2 <- getNode(tree, node2)
25
26
        # most recent common ancestor
27
         comAnc <- MRCA(tree, t1, t2)
28
         desComAnc <- descendants (tree, comAnc)
29
        # path: common ancestor to t1
30
31
         ancT1 <- ancestors(x, t1)
32
         path1 <- intersect (desComAnc, ancT1)
33
        # path: common ancestor to t2
34
         ancT2 \leftarrow ancestors(x, t2)
35
36
         path2 <- intersect (desComAnc, ancT2)
37
38
        # union of the path above paths
39
         retval <- union(path1, path2)
40
```

2.2 RSnow

The RSnow implementaion builds on top of original R version by parallelizing the descendants function. In order to make independent subproblems, for a given node, every other node keeps marching upwards to its ancestors until they either reached the root or encountered the given node. We were unable to parallelize the ancestor function. The original R version seems to have taken the most efficent serial approach for calculating a given node's ancestors.

Pseudocode: (see code in Appendix A.1)

```
descendants(tree, given_node){
    #let x be the list of all nodes except the given_node
    for i from 1 to height of tree
        if x == given_node -> append node to retval
```

```
5 update x with its correponding ancestor
6 return retval
7 }
```

2.3 OpenMP

ALICIA WRITE STUFF HERE

See code in Appendix A.2.

2.4 CUDA

Our CUDA implementation of the shortestPath function utilizes the GPU to find all ancestors of a given pair of nodes and then construct the shortest path between them. Our implementation assumes CSIF's pc43's resources, which are 1024 threads per block and 1 GB of global memory. We assume the given data can fit in our GPUs global memory. This assumption may limit the test we will be able to perform. Our solution utilizes the fact that the shortest path between two nodes in a tree must converge at the lowest common ancestor of both nodes. In cases, where one node is an ancestor of another, then the shortest path is then found by traversing the parents of the child node. We parallelized our code by finding both sets of ancestors of the given nodes at the same time. Since neither node needs to know about the other to find its own ancestors, this problem can be done independently of each other. Both sets of ancestors are then traversed to find the shortest path. We were unable to parallelize this part of the solution since each list of ancestors must be checked to find overlapping elements.

See code in Appendix A.3.

3 Experiment Results

These are the results of running the above scripts with a simplified internet (i.e. n = 6).

```
>> i = [ 2 6 3 4 4 5 6 1 1];
>> j = [ 1 1 2 2 3 3 3 4 6];
>> n = 6;
>> G = sparse(i,j,1,n,n);
>> Finaltimetest
```

pagerank1	pagerank2	pagerank3A	pagerank3B	pagerankpow
0.0002	0.0001	0.0002	0.0003	0.0004

These are the results of running the above scripts with the Harvard 500 dataset (i.e. n = 500).

- >> load Harvard500
- >> Finaltimetest

pagerank1	pagerank2	pagerank3A	pagerank3B	pagerankpow
0.0024	0.0329	0.0226	0.0011	0.0255

These results are consistent with the dicussion above.

4 Discussion

4.1 RSnow

TALK ABOUT BIG O's.

4.2 OpenMP

TALK ABOUT BIG O's.

4.3 CUDA

THIS IS JUST A GUESS.

Compared to the serial version, the cuda implementation performed slower in most test cases. While the cuda version can compute both given nodes ancestors at the same time, it must also load the entire tree into the GPUs memory.

5 Conclusion

The PageRank algorithm was the starting point of Google's rise to fame. It was able to numerically quantify the "quality" on links/web pages of the internet. Pagerank is a Markov chain for which we solve for the dominant eigenvector of its transition probability matrix. There are two main methods of solving such a system of linear equations. The Power method is shown here to be the best method because of its efficiency in run time and memory usage. The run times are decent, as Hopcraft stated, it varies logarithmically with the size of input (n web paes). For the current day, the sheer amount of data that needs to processed is daunting. Further studies on PageRank could be done in further optimizing its space usage.

6 Acknowledgements

We would like to thank Professor Norman Matloff for his guidance and knowledge presented during lectures. This work is the result of a final project for ECS 158 Winter Quarter 2015. His open-source textbook, blog and various tutorial were an essential part of our learning. We would like to thank the teaching assistance Shengren Li for offering invaluble advice and feedback on our codes (especically our CUDA) throughout the quarter.

7 Appendix

A Codes

INSERT ALL CODES HERE ALONG WITH A PARAGRAPH EXPLAINING IT

A.1 RSnow Code

```
1
   SNOW <- function(x, size , root , type=c("descendants")){</pre>
 2
        ans \leftarrow \text{rep}(0, \text{size})
 3
        mystart \leftarrow (myid-1)*length(x)+1
        myend \leftarrow myid*length(x)
 4
 5
 6
        type <- match.arg(type)
 7
         if (type == "descendants"){
 8
             v1 <- descendant
 9
             v2 <- ancestor
10
             #initalization
11
             temp <- v1 [mystart:myend]
12
             #second and beyond iteration
13
14
             for (j in 1:size){
                  if (node %in% temp){
15
                       setthese <- which (temp == node) + mystart-1
16
17
                       ans[setthese] <- 1
18
                  blah \leftarrow \operatorname{rep}(-1, \operatorname{length}(\operatorname{temp}))
19
20
                  for (i in (1:length(temp))){
21
                       matched_pos \leftarrow which(v1 = temp[i])
22
                       if (length (matched_pos) != 0){
23
                           blah [which (temp = temp[i])] <- matched_pos
24
25
                       else {\#matched_pos == 0
                           ## R is 1 INDEXED!
26
                           if (type = "descendants") {
27
28
                                blah [ i ] <- 1
29
                       }
30
                  }#for i
31
32
                  #"go to your parents set"
                  difference <- length(temp) - length(v2[blah])
33
34
                  temp \leftarrow v2[blah]
                  if (difference > 0){
35
36
                      temp <- c(rep(0, difference), temp)
38
                  if (node %in% temp) {
39
                       setthese \leftarrow which (temp = node) + mystart-1
                       ans[setthese] <- 1
40
41
             }#j loop
42
        }#new endif for type=descendants
43
        return (ans)
44
    }# end SNOW
45
46
    setmyid <- function(i){</pre>
47
        myid <<-i
48
49
   }
50
51
   ## get descendants with RSnow
   RSnowdescendants <- function (phy, node, type=c("tips", "children", "all"), cls) {
52
        type <- match.arg(type)
53
54
        ## look up nodes, warning about and excluding invalid nodes
        oNode <- node
56
        node <- getNode(phy, node, missing="warn")</pre>
57
```

```
is Valid <- !is.na(node)
58
59
        node <- as.integer(node[isValid])
60
        if (type == "children") {
 61
            res <- lapply (node, function (x) children (phy, x))
63
            ## if just a single node, return as a single vector
            if (length(res)==1) res \leftarrow res [[1]]
64
65
        } else {
            ## edge matrix must be in preorder for the C function!
66
            #if (phy@order=="preorder") {
67
 68
                edge <- phy@edge
 69
            #} else {
 70
            #
                  edge <- reorder(phy, order="postorder")@edge
            #}
 71
            ## extract edge columns
 72
            ancestor <- as.integer(edge[, 1])
 73
 74
            descendant <- as.integer(edge[, 2])
 75
            ## return indicator matrix of ALL descendants (including self)
 76
 77
            #isDes <- .Call("descendants", node, ancestor, descendant)</pre>
            clusterExport(cls,c("node", "ancestor", "descendant", "setmyid", "SNOW"),
 78
                envir=environment())
            dexgrps <- splitIndices(length(ancestor), length(cls))
 79
 80
            rootdex \leftarrow which(phy@edge[,1] == 0)
            clusterApply(cls,1:length(cls),setmyid)
 81
            newisDes <- clusterApply(cls,dexgrps,SNOW,length(ancestor),rootdex,</pre>
 82
                descendants")
            isDes <- (matrix(Reduce('+',newisDes),nrow=length(ancestor),ncol=1))
 83
            storage.mode(isDes) <- "logical"</pre>
 84
 85
            ## for internal nodes only, drop self (not sure why this rule?)
            int.node <- intersect (node, nodeId (phy, "internal"))
 86
            isDes[cbind(match(int.node, descendant),
 87
 88
                match(int.node, node)) < FALSE
 89
90
            ## if only tips desired, drop internal nodes
            if (type="tips") {
91
                isDes [descendant %in% nodeId(phy, "internal"),] <- FALSE
92
93
            ## res <- lapply(seq_along(node), function(n) getNode(phy,
94
95
                    descendant [isDes[,n]]))
96
            res <- getNode(phy, descendant[isDes[, seq_along(node)]])
97
        ## names(res) <- as.character(oNode[isValid])
98
99
100
        res
101
    }
102
103
    104
    # shortestPath
105
    106
    RSnowshortestPath <- function(phy, node1, node2, cls){
107
108
109
        ## conversion from phylo, phylo4 and phylo4d
110
        if (class(phy) = "phylo4d") {
111
            x <- extractTree(phy)
112
        else if (class(phy) != "phylo4"){
113
114
            x \leftarrow as(phy, "phylo4")
```

```
115
116
        ## some checks
117
        t1 \leftarrow getNode(x, node1)
118
        t2 \leftarrow getNode(x, node2)
         if (any (is.na(c(t1,t2)))) stop("wrong node specified")
119
120
         if (t1==t2) return (NULL)
121
122
        ## main computations
        comAnc \leftarrow MRCA(x, t1, t2) \# common ancestor
123
124
        desComAnc <- RSnowdescendants(x, comAnc, type="all", cls)
        ancT1 <- ancestors(x, t1, type="all")
125
126
        path1 <- intersect (desComAnc, ancT1) # path: common anc -> t1
127
128
        ancT2 <- ancestors(x, t2, type="all")
129
        path2 <- intersect (desComAnc, ancT2) # path: common anc -> t2
130
131
        res <- union(path1, path2) # union of the path
132
        ## add the common ancestor if it differs from t1 or t2
        if (!comAnc %in% c(t1,t2)) {
133
134
             res <- c(comAnc, res)
135
136
137
         res <- getNode(x, res)
138
139
        return (res)
140
    } # end shortestPath
```

A.2 OpenMP Code

```
ALICIA CODE GOES HERE
```

A.3 CUDA Code

```
#include <stdio.h>
    #include <stdlib.h>
    #include <cuda.h>
 5
    //nvcc shortestPath.cu
 7
    struct node
 8
 9
          int nodeID, ancestor;
10
          char label [20]; //max size of label is 20
11
12
     {\tt void} \  \, {\tt setNode} ( {\tt node} \ \& {\tt phy} \, , \  \, {\tt int} \  \, {\tt numNodes} \, , \  \, {\tt int} \  \, {\tt id} \, , \  \, {\tt int} \  \, {\tt aID} \, , \  \, {\tt char} \  \, * \  \, {\tt label} ) 
13
14
    {
          phy.nodeID = id;
15
         phy.ancestor = aID;
memset(phy.label, '\0', sizeof(label));
strcpy(phy.label, label);
16
17
18
19 }//setNode
20
    __global__ void kernel(node * array, int numNodes, int id1, int id2,
22
                                       int * ancestorID1, int * ancestorID2)
23 {
```

```
24
       int idx = blockIdx.x * blockDim.x + threadIdx.x;
25
       if (idx < numNodes)</pre>
26
27
            if (array[idx].nodeID == id1) //if found target node
28
29
                int ancestorIndex = 0;
30
                \verb|node| temp| = \verb|array|[idx]|; //start| from current| node
                while (temp.ancestor != 0)
31
32
                     ancestorID1[ancestorIndex++] = temp.ancestor; //add all ancestors
33
34
                     for (int i=0; i<numNodes; i++)
35
                         if (array[i].nodeID == temp.ancestor)
36
37
38
                              temp = array[i];
39
                              break;
40
                         }//if
                    }//for
41
42
                }//while
            }//if
43
44
            else if (array[idx].nodeID == id2) //if found target node
45
46
                int ancestorIndex = 0;
47
                node temp = array[idx]; //start from current node
                while (temp.ancestor != 0)
48
49
                     ancestorID2[ancestorIndex++] = temp.ancestor; //add all ancestors
50
51
                     for (int i=0; i<numNodes; i++)
52
53
                          if (array[i].nodeID == temp.ancestor)
54
                         {
55
                              temp = array[i];
56
                              break;
57
                    }//for
58
           }//while
}//if
59
60
61
       }//if
62
   }//kernel
63
64
   void shortestPath(node * phy, int numNodes, char * label1, char * label2)
65
66
67
       node * deviceArray;
68
       int * deviceID1;
69
       int * deviceID2;
70
       int * ancestorID1 = new int[numNodes]; //initialize max size to number of nodes
71
       int * ancestorID2 = new int[numNodes];
72
       float blockSize = 1024; //num threads per block
73
74
       //check if invalid query
75
       node temp1, temp2;
76
       for (int i=0; i<numNodes; i++)
77
            ancestorID1[i] = 0;
78
            ancestorID2[i] = 0;
79
80
             if \ (\mathtt{strcmp}(\mathtt{label1}\,,\ \mathtt{phy}[\mathtt{i}\,].\,\mathtt{label}) =\!\!= 0) 
81
                temp1 = phy[i];
82
            else if (strcmp(label2, phy[i].label) == 0)
83
                temp2 = phy[i];
84
85
       if ((temp1.ancestor == temp2.nodeID) || (temp2.ancestor == temp1.nodeID))
86
87
88
            printf("named integer(0)\n");
89
            return;
90
       }//if
91
```

```
//allocate device memory
 92
 93
         \verb|cudaMalloc(\&deviceArray|, sizeof(node) * numNodes);|
         cudaMalloc(&deviceID1, sizeof(int) * numNodes);
 94
 95
         cudaMalloc(&deviceID2, sizeof(int) * numNodes);
 96
         cudaMemcpy(deviceArray, phy, sizeof(node) * numNodes, cudaMemcpyHostToDevice);
 97
         \verb|cudaMemcpy| (\verb|deviceID1|, ancestorID1|, size of (int) * numNodes, cudaMemcpyHostToDevice); \\
         \verb|cudaMemcpy| (\texttt{deviceID2}\,, \texttt{ ancestorID2}\,, \texttt{ sizeof}(\texttt{int}) * \texttt{numNodes}\,, \texttt{ cudaMemcpyHostToDevice});\\
 98
 99
100
        dim3 dimBlock(blockSize);
        dim3 dimGrid(ceil(numNodes/blockSize));
101
102
103
         //compute ancestors
        kernel <<< dimGrid, dimBlock >>> (deviceArray, numNodes, temp1.nodeID, temp2.nodeID, ←
104
             deviceID1, deviceID2);
105
         \verb|cudaMemcpy| (ancestorID1 \;, \; deviceID1 \;, \; sizeof(int) \; * \; numNodes \;, \; cudaMemcpyDeviceToHost); \\
106
         cudaMemcpy(ancestorID2, deviceID2, sizeof(int) * numNodes, cudaMemcpyDeviceToHost);
107
         cudaFree(deviceArray);
108
         cudaFree(deviceID1);
         cudaFree(deviceID2);
109
110
111
         // {
m find} shortest path
112
         int * path = new int[numNodes];
113
        int currentPath = ancestorID1[0];
114
         int pathIndex = 0;
         bool isLCAPath = false;
115
116
         //check if path converges at LCA
117
         for (int i=0; i<numNodes; i++)
118
119
             path[i] = 0;
120
             if (temp1.nodeID == ancestorID2[i])
121
122
                  for (int j=0; j<i; j++)
                      path[j] = ancestorID2[j];
123
124
                  isLCAPath = true;
                  break;
125
126
             }//if
127
             else if (temp2.nodeID == ancestorID1[i])
128
129
                  for (int j=0; j<i; j++)
130
                       path[j] = ancestorID1[j];
131
                  isLCAPath = true;
132
                  break:
133
             }//else if
        }//for
134
135
136
         //if one node is the ancestor of another
137
         if (!isLCAPath)
138
139
              for(int i=0; i<numNodes; i++)
140
141
                  for (int j=0; j<numNodes; j++)
142
143
                       if (currentPath == ancestorID2[j])
144
                           break;
145
                       if ((ancestorID2[j] = 0) || (j = numNodes-1))
146
                           {\tt path}\,[\,{\tt pathIndex} + +] \,=\, {\tt ancestorID1}\,[\,{\tt i}\,]\,;
147
148
                           currentPath = ancestorID1[i];
149
                           break;
150
                       }//if
                  }//for
151
             }//for
152
153
154
             if (pathIndex == 0)
155
                  path[pathIndex++] = currentPath;
156
157
             for (int i=0; i<numNodes; i++)
158
```

```
if (ancestorID2[i] = currentPath)
159
160
                         break;
                    \mathtt{path}\,[\,\mathtt{pathIndex} + +] \,=\, \mathtt{ancestorID2}\,[\,\mathtt{i}\,]\,;
161
162
               }//for
          }//if
163
          for (int i=0; i<numNodes; i++)
164
165
166
               if (path[i] == 0)
                    break;
167
               for (int j=0; j<numNodes; j++)
168
169
170
                    if (path[i] == phy[j].nodeID)
171
                    {
172
                         {\tt printf\left(``\%s\ ''\ ,\ phy\left[\,j\,\right].\,label\,\right);}
173
                         break;
                    }//if
174
         }//for
}//for
175
176
177
         printf("\n");
178
         for (int i=0; i<numNodes; i++)
179
180
181
               if (path[i] == 0)
182
                    break;
               printf("%d ", path[i]);
183
         }//for
184
         printf("\n");
185
186
187
          delete [] ancestorID1;
                   [] ances
[] path;
188
          delete
                      ancestorID2;
189
          delete
190
    }//shortestPath
```

B Who Did What

Alicia wrote the OpenMP implementation. Bryan wrote the CUDA implementation. Raymond wrote the RSnow implementation. We worked on running tests and writing the report in LATEX.

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