

Relational Algebra

Relational algebra is a **procedural query language** which works on relational models. Procedural query language tells what data to be retrieved and how to be retrieved.

A fundamental property is that every operator in the algebra accepts (one or two) relation instances as arguments and returns a relation instance as the result.

A relational algebra expression is recursively defined to be a relation, a unary algebra operator applied to a single expression, or a binary algebra operator applied to two expressions.

The fundamental operations in the relational algebra are **select**, **project**, **union**, **set difference**, **Cartesian product**, **and rename**.

The select, project, and rename operations are called <u>unary operations</u>, because they operate on one relation.

The other three operations operate on pairs of relations and are, therefore, called binary operations.

Set Operation

- 1. Union
- 2. Intersection
- 3. Set difference
- 4. Cartesian product

Union

- Union of two relations R and S (R U S) defines a relation that contains all the tuples of R, or S, or both R and S, duplicate tuples being eliminated.
- R and S must be union-compatible.
- If R and S have I and J tuples, respectively, union is obtained by concatenating them into one relation with a maximum of (I + J) tuples.
- UNION is symbolized by U symbol.



Example:

Graduates

Number	Surname	Age
7274	Robinson	37
7432	O'Malley	39
9824	Darkes	38

Managers

managere .			
Number	Surname	Age	
9297	O'Malley	56	
7432	O'Malley	39	
9824	Darkes	38	

Graduates \cup **Managers**

Number	Surname	Age	
7274	Robinson	37	
7432	O'Malley	39	
9824	Darkes	38	
9297	O'Malley	56	

Intersection

- The intersection of two relations R and $S(R \cap S)$, defines a relation consisting of the set of all tuples that are in both R and S.
- R and S must be union-compatible.
- Represented using basic operations: $R \cap S = R (R S)$

Example:

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Number	Surname	Age
7274	Robinson	37
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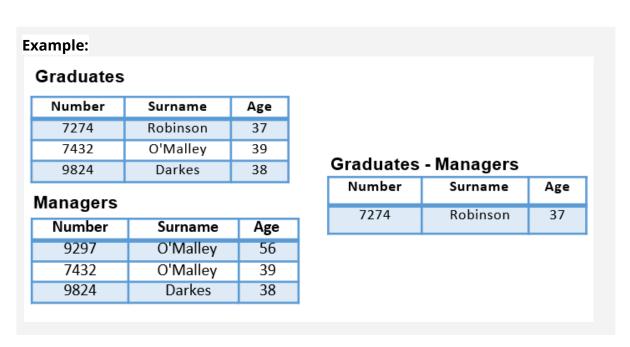
Graduates ∩ **Managers**

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Set Difference

- The difference of two relations R and S (R S), define a relation consisting of the tuples that are in relation R, but not in S.
- R and S must be union-compatible.
- It is denoted by(-).
- Represented using basic operations: R -S



Cartesian product

- The Cartesian product of two relations R and S (R X S), defines a relation between every tuple of relation R with every tuple of relation S.
- It is denoted by (X).
- Represented using basic operations: R X S



Example:

Student

Sr No	Name	Grade
1	Andrew	D
2	Robin	В

Data

Seat No	Age
18	25
11	21

Student x Data

Sr No	Name	Grade	Seat No	Age
1	Andrew	D	18	25
1	Andrew	D	11	21
2	Robin	В	18	25
2	Robin	В	11	21

Unary Relational operations

- 1. Select
- 2. Project
- 3. Rename

Select Operation

- The select operation is performed to select certain rows or tuples of a table, so it performs its action on the table horizontally.,
- The tuples are selected through this operation using a predicate or condition.
- It works on a single table and takes rows that meet a specified condition, copying them into a new table.
- Denoted by lower Greek letter sigma (σ).

Relation: σ predicate(R)



Example of Selection:

Employees

Surname	FirstName	Age	Salary
Smith	Mary	25	2000
Black	Lucy	40	3000
Verdi	Nico	36	4500
Smith	Mark	40	3900

σ _{Age<30 v Salary>4000} (Employees)

Surname	FirstName	Age	Salary
Smith	Mary	25	2000
Verdi	Nico	36	4500

In selection operation the comparison operators like <, >, =, <=, >=, <> can be used in the predicate

Project Operation

- The Select operation operates horizontally on the table. Conversely, the Project operator works on a single table vertically.
- It is a unary operation that returns a relation that includes a **subset of the attributes** of the operand.
- Since the relation is a set, any duplicate rows are eliminated.
- Projection is represented by a Greek letter (Π).



Example of Projection:

Employees

Surname	FirstName	Department	Head
Smith	Mary	Sales	De Rossi
Black	Lucy	Sales	De Rossi
Verdi	Mary	Personnel	Fox
Smith	Mark	Personnel	Fox

$\pi_{Surname, FirstName}$ (Employees)

Surname	FirstName	
Smith	Mary	
Black	Lucy	
Verdi	Mary	
Smith	Mark	

Rename

- This is a unary operator which changes attribute names for a relation without changing any values.
- Renaming removes the limitations associated with set operators.
- Representation: ρ old name(R) \rightarrow New Name(R)



Example: ρ Father->Parent(Paternity)

Paternity

Father	Child	
Adam	Cain	
Adam	Abel	
Abraham	Isaac	
Abraham	Ishmael	

$\rho_{\text{Father} -> \, \text{Parent}}(\text{Paternity})$

Parent	Child	
Adam	Cain	
Adam	Abel	
Abraham	Isaac	
Abraham	Ishmael	

Join Operation

- The JOIN operation, denoted by ⋈, is used to combine related tuples from two relations into single "longer" tuples.
- The join operator allows the combination of two relations to form a single new relation.
- The JOIN operation can be specified as a CARTESIAN PRODUCT operation followed by a SELECT operation.

Natural Join

An Equijoin of the two relations R and S over all common attributes x. One occurrence of each common attribute is eliminated from the result.

Hence the degree is the sum of the degrees of the relations R and S less the number of attributes in x



$\mathbf{r_1}$

<u>Employee</u>	Department	
Smith	sales	
Black	production	
White	production	

r_2

Department	Head
production	Mori
sales	Brown

r1 ⋈r2

Employee	Department	Head
Smith	sales	Brown
Black	production	Mori
White	production	Mori

Theta Join

Defines a relation that contains tuples satisfying the predicate F from the Cartesian product of R and S.

The predicate F is of the form R.ai θ S.bi where θ may be one of the comparison operators (<, \leq , >, \geq , =, \neq).