#### **Project Report**

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## **Project Statement**

Study and Implementation of Distributed SSH Key Management with Proactive RSA Threshold Signatures

### Abstract

Large enterprises often face difficulty in managing the high number of SSH keys. In this report, we study and implement ESKM - a distributed enterprise SSH key manager [1]. ESKM is a secure and fault-tolerant logically-centralized SSH key manager. In ESKM architecture, SSH private keys are never stored at any single node. Instead, each node only stores a share of the private key. These shares are refreshed at regular intervals for enforced security. For signing, the system uses k-out-of-n threshold RSA signatures.

### **Notations**

 $\bullet$  k: threshold value

•  $\ell$ : number of nodes

• n: RSA public key modulus

 $\bullet$  e: RSA public exponent

 $\bullet$  d: RSA private exponent

• p, q: safe primes

 $\bullet \ \Delta : \ell!$ 

•  $s_i$ : current share of node i

•  $x_i$ : signature fragment by node i

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• H: digest to be signed

•  $\sigma$ : final signature

• sk : RSA private key

• pk : RSA public key

• dk : RSA dummy private key

# 1 The Security Manager

• KeyGen: Takes  $b \in \{1024, 2048, 4096\}$  as input and proceeds as follows:

- Generate  $\mathsf{sk}$  and  $\mathsf{dk}$  b-size RSA key:

openssl genrsa -out sk.pem [b]

- Get pk from sk:

openssl rsa -in sk.pem -outform PEM -pubout -out pk.pem

- Convert pk into pkssh:

ssh-keygen -f pk.pem -i -mPKCS8

- Extract e, n, p, q from sk with regex

- Set  $p' = \frac{p-1}{2}, q' = \frac{q-1}{2}$ 

- Set m = p'q'
- $\text{ Set } d = e^{-1} \mod m$
- Output pkssh, dk
- $\bullet$  SplitSecret: Takes d as input and proceeds as follows:
  - -Set  $a_0 = d$
  - Choose  $a_i \in_R \{0, 1, ..., m 1\}$  ∀ 1 ≤  $i \le k 1$ .
  - Set  $f(x) = \sum_{i=0}^{k-1} a_i x^i$
  - Compute  $s_i = f(i) \mod m, \forall \ 1 \le i \le \ell$
  - Output  $s_i$ ,  $\forall 1 \leq i \leq \ell$

Note: All group operations from here will be done mod n unless mentioned otherwise.

- GenVerifyInfo: Takes p, q as input and proceeds as follows:
  - Set  $g_p$  = Primirive Root of p
  - Set  $g_q$  = Primitive Root of q
  - Use the Chinese Remainder Theorem to get the primitive root of n:
    - $* \bar{g} = g_p \mod p$
    - $* \bar{g} = g_q \mod q$
  - $Set g = \bar{g}^2$
  - Output  $q^{a_i}$ ,  $\forall 0 < i < k-1$
- Broadcast: Proceeds as follows:
  - Send pkssh, dk to client
  - $\forall 1 \le i \le \ell$ , send  $(g^{a_j}, \forall 0 \le j \le k-1), s_i$  to node i
  - Delete sk

### 2 The Control Cluster

- VerifyShare: Takes  $s_i$  as input and proceeds as follows:
  - Verify  $g^{s(i)} = \prod_{j=0}^{k-1} (g^{a_j})^{i^j}$
- $\bullet$  Sign: Takes H as input and proceeds as follows:
  - Compute  $x_i = H^{2\Delta s(i)}$
  - Send  $x_i$  to client
- ullet RefreshShare: Takes  $s_i$  as input and proceeds as follows:
  - $\text{ Set } a_0 = 0$
  - Choose  $a_i \in_R \{0, 1, ..., n-1\} \ \forall \ 1 \le i \le k-1$ .
  - Set  $z(x) = \sum_{i=0}^{k-1} a_i x^i$

- $\forall 1 \le i \le \ell$ , send z(i) to node i
- $\ \forall \ 1 \leq i \leq \ell,$ receive $s_i'$  from node i
- If  $|\{s'_1, s'_2, ..., s'_\ell\}| \ge k$ 
  - \* Compute  $s_i^* = s_i + \sum s_i' \in \mathbb{Z}$
  - \* Set  $s_i = s_i^*$
- Repeat Refresh Share after every 60 seconds '  $\,$

### 3 The Client

- CombineSig: Takes  $x_i$  from at least k servers as input and proceeds as follows:
  - Compute  $w = \prod_i x_i^{2\lambda_i^s}$  where,  $\lambda_i^s = \Delta \lambda_i \in \mathbb{Z}$
  - Compute  $a = (4\Delta^2)^{-1} \mod e$
  - Compute  $b = \frac{(1 (4\Delta^2 a))}{e} \in \mathbb{Z}$
  - Compute  $\sigma = w^a \times H^b$
  - Output  $\sigma$

# 4 Algorithms Implemented

- Square and Multiply: exponentiation using pow() Note: If  $c < 0, x^c \equiv (x^{-1})^{|c|}$
- Extended Euclidean Algorithm: modular multiplicative inverse
- Primitve Root of (safe prime) Note: x is a primitive root if  $\forall$  prime factors  $p_i$  of  $(p-1), x^{(p-1)/p_i} \not\equiv 1 \mod p$
- $\bullet$  Chinese Remainder Theorem: primitive root of n
- Horner's Method: polynomial evaluation

### 5 Miscellaneous

- Patched OpenSSL:
  - Default install location: /usr/local/ssl/bin
  - Generates sk with p and q as safe primes
  - Generates dk with random values for all fields except for e and n.
- Working directory: /tmp
- All communications are over TLS

#### 6 Source Code

- Core
- OpenSSL

# References

- [1] Y. Harchol, I. Abraham, and B. Pinkas. Distributed ssh key management with proactive rsa threshold signatures. Cryptology ePrint Archive, Report 2018/389, 2018. https://eprint.iacr.org/2018/389
- [2] Tatu Ylonen: SSH Secure Login Connections over the Internet. Proceedings of the 6th USENIX Security Symposium, pp. 37-42, USENIX, 1996.