COMPUTER SCIENCE E-20, SPRING 2014

Homework Problems Induction I, Strong Induction

Due Thursday, February 12, 2015 before 9PM EST. Upload a PDF of your answers at https://canvas.harvard.edu/courses/1815/assignments/17757

1. Prove that for all nonnegative integers n

$$\sum_{i=0}^{n} i^3 = \left(\sum_{i=0}^{n} i\right)^2$$

Hint: the following identity may be useful

$$\sum_{i=0}^{n} i = \frac{n(n+1)}{2}$$

Solution:

 \bullet Proof: for all nonnegative integers n

$$\sum_{i=0}^{n} i^{3} = \left(\sum_{i=0}^{n} i\right)^{2}$$

• The Induction hypothesis P(n), is:

$$0^{3} + 1^{3} + 2^{3} \dots + n^{3} = (0 + 1 + 2 + \dots + n)^{2}$$
$$0^{3} + 1^{3} + 2^{3} \dots + n^{3} = (n(n+1)/2)^{2}$$

- Base Case (n = 0) : 0 == 0
- Base Case (n = 1) : 1 == 1
- Inductive Step: Assume P(n) is true $\forall n \geq 0$, and prove P(n+1)

$$0^{3} + 1^{3} + 2^{3} \dots n^{3} + (n+1)^{3} = (0+1+2+\dots+n+n+1)^{2}$$
$$0^{3} + 1^{3} + 2^{3} \dots n^{3} + (n+1)^{3} = ((n+1)(n+2)/2)^{2}$$

from Inductive Step:

$$0^3 + 1^3 + 2^3 \dots n^3 = (n(n+1)/2)^2$$

Therefore

$$(n(n+1)/2)^{2} + (n+1)^{3} = ((n+1)(n+2)/2)^{2}$$

$$(n+1)^{2}/2^{2} (n^{2} + 4(n+1)) = ((n+1)(n+2)/2)^{2}$$
$$(n+1)^{2}/2^{2} (n^{2} + 4n + 4)) = ((n+1)(n+2)/2)^{2}$$
$$(n+1)^{2}/2^{2} (n+2)^{2}) = ((n+1)(n+2)/2)^{2}$$
$$(n+1)^{2}(n+2)^{2}/2^{2}) = ((n+1)(n+2)/2)^{2}$$
$$(n+1)(n+2)/2)^{2}) = ((n+1)(n+2)/2)^{2}$$

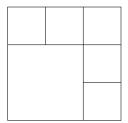
This proves P(n+1), completing the proof by induction.

- 2. Consider the sequence $a_1 = 1, a_2 = 3, ..., a_n = a_{n-1} + a_{n-2}$. Using strong induction prove that $a_n \leq \left(\frac{7}{4}\right)^n$ for all positive integers n.
 - Solution:
 - Proof: By strong induction on n that for all positive integers $n, a_n \leq \left(\frac{7}{4}\right)^n$.
 - The Induction hypothesis P(n), is: $a_n \leq \left(\frac{7}{4}\right)^n \forall n \geq 1$
 - Base Case (n = 1): $a_1 \le \left(\frac{7}{4}\right)^1 = 1 \le \left(\frac{7}{4}\right) = 1 \le 1.75$
 - Base Case (n = 2): $a_2 \le \left(\frac{7}{4}\right)^2 = 3 \le \left(\frac{7}{4}\right)^2 = 3 \le 1.75 * 1.75 = 3 \le 3.06$
 - Inductive Step : Assume P(n) is true \forall n \geq 1, and prove P(n+1) : $a_{n+1} \leq {7 \choose 4}^{n+1}$ substituting (n+1) in $a_n = a_{n-1} + a_{n-2}$ $a_{n+1} = a_{n+1-1} + a_{n+1-2} \leq {7 \choose 4}^{n+1}$ $= a_n + a_{n-1} \leq {7 \choose 4}^{n+1}$ $= {7 \choose 4}^n + {7 \choose 4}^{n-1} \leq {7 \choose 4}^{n+1}$ $= {7 \choose 4}^{n-1} * {7 \choose 4}^{n-1} \leq {7 \choose 4}^{n+1}$ $= {7 \choose 4}^{n-1} * {7 \choose 4}^{n} + {1 \choose 4}^{n} \leq {7 \choose 4}^{n} * {7 \choose 4}^{n}$ $= {7 \choose 4}^{n-1} * {7 \choose 4}^{n-$

This proves P(n+1), completing the proof by induction.

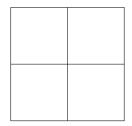
3. Prove using strong induction that any square can be subdivided into n smaller squares, where n > 5. For example, the large square below has been subdivided into 6 squares.

Hint: first show that any square subdivided into k squares can easily be subdivided into k+3 squares, then think how many base cases you need show are true (it is not just the case of n=6).



Solution:

- Proof: By strong induction on n that any square can be subdivided into n smaller squares, where n > 5.
- The simplest way a square can be subdivided into smaller squares is by dividing it into 4 equal squares as shown in the picture below.

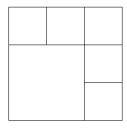


Each of these smaller squares represent a square, and so can again be divided in the same way into 4 smaller squars, as shown below.

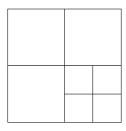


This shows that any square subdivided into k squares can always be divided into K+3 squares.

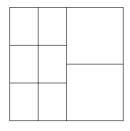
• Base Case (n = 6):



• Base Case (n = 7):



• Base Case (n = 8):



• Inductive Step: Assume P(n) is true for all n > 5, and we now prove P(n+1):

when P(n+1)=6,7,8 , these form our base cases and have been proved above

when P(n+1) > 8: From the inductive step we know that any square can be divided into P(n) smaller squares. We have also shown above that a single square k can always be divided into k+3 squares. Thus P(n+1) can be divided into P(n)+4 smaller squares. This completes the prove by induction.

4. The Fibonacci numbers are defined by $F_0 = 0, F_1 = 1, F_n = F_{n-1} + F_{n-2}$. Prove using strong induction that for all $n \ge 0, F_{3n}$ is even.

Solution:

- Proof: By strong induction on n that for all $n \geq 0$, F_{3n} is even.
- The Induction hypothesis P(n), is: F_{3n} is even $\forall n \geq 0$
- Base Case (n = 1): $F_3 = F_2 + F_1 = F_1 + F_0 + F_1 = 1 + 0 + 1 = 2$.
- $\begin{array}{l} \bullet \ \, \mathrm{Base} \ \, \mathrm{Case} \ \, (n=2) : \, F_6 = F_5 + F_4 \\ F_4 = F_3 + F_2 = 2{+}1 = 3 \\ F_6 = F_4 + F_3 + F_4 = 3 + 2 + 3 = 8 \end{array}$
- Inductive Step : Assume P(n) is true \forall n \geq 0, and prove P(n+1) : F_{3(n+1)} is even substituting (n+1) in $F_n = F_{n-1} + F_{n-2}$

$$F_{3(n+1)} = F_{n+1-1} + F_{n+1-2}$$

$$F_{3(n+1)} = F_n + F_{n-1}$$

From the inductive hypothesis P(n): F_n is even. So, $F_n = 2*a(someintegera)$ F_{n-1} is even. So, $F_{n-1} = 2*b(someintegerb)$

$$F_{3(n+1)} = 2 * a + 2 * b$$

 $F_{3(n+1)} = 2(a+b)$, which implies $F_{3(n+1)}$ is even.

This proves P(n+1), completing the proof by induction.