

# A Type Theory with Native Homotopy Universes

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We present a type theory called  $\lambda \simeq_2$  with an extensional equality relation; that is, the universe of types is closed by reflection into it of the logical relation defined by induction on the structure of types.

The type system has three universes:

- The universe **Prop** of *propositions*. An object of **Prop** is called a *proposition*, and the objects of a proposition are called *proofs*. There is a proposition **1**, which has a unique proof  $*$ .
- The universe **Set** of *sets*.
- The universe **Groupoid** of *groupoids*.

The system has been designed in such a way that it should be straightforward to extend the system with three, four, ... dimensions.

For each universe  $U$ , we have an associated relation of equality between  $U$ -types  $\simeq$ , and between objects of  $U$ -equal types  $\sim$ . The associated rules of deduction are:

$$\frac{A : U \quad B : U}{A \simeq B : U} \quad \frac{a : A \quad e : A \simeq B \quad b : B}{a \sim_e b : U^-}$$

where  $U^-$  is the universe one dimension below  $U$ . Thus:

- Given two propositions  $\phi$  and  $\psi$ , we have the proposition  $\phi \simeq \psi$  which denotes the proposition ‘ $\phi$  if and only if  $\psi$ ’. If  $\delta : \phi$ ,  $\epsilon : \psi$  and  $\chi : \phi \simeq \psi$ , then  $\delta \sim_\chi \epsilon = \mathbf{1}$ . (Cf In homotopy type theory, any two objects of a proposition are equal.)
- Given two sets  $A$  and  $B$ , we have the set  $A \simeq B$ , which denotes the set of all bijections between  $A$  and  $B$ . Given  $a : A$ ,  $f : A \simeq B$  and  $b : B$ , we have the proposition  $a \sim_f b : \mathbf{Prop}$ , which denotes that  $a$  is mapped to  $b$  by the bijection  $f$ .
- Given two groupoids  $G$  and  $H$ , we have the groupoid  $G \simeq H$ , which denotes the groupoid of all groupoid isomorphisms between  $G$  and  $H$ . Given  $g : G$ ,  $\phi : G \simeq H$  and  $h : H$ , we have the set  $g \sim_\phi h : \mathbf{Set}$ , which can be thought of as the set of all paths between  $\phi(g)$  and  $h$  in  $H$ .

The introduction and elimination rules for  $\simeq$  ensure that  $A \simeq B$  is the type of equivalences between  $A$  and  $B$ .

$$\begin{array}{c} \frac{A : U}{1_A : A \simeq A} \quad \frac{a : A}{r_a : a \sim_{1_A} a} \quad \frac{e : A \simeq B \quad a : A}{e^+(a) : B} \quad \frac{e : A \simeq B \quad b : B}{e^-(b) : A} \\[10pt] \frac{a : A \quad b : B \quad e : A \simeq B}{e^=(a, b) : (a \sim_{1_A} e^-(b)) \simeq (e^+(a) \sim_{1_B} b)} \quad \frac{\begin{array}{l} \Gamma, x : A \vdash b : B \\ \Gamma, y : B \vdash a : A \\ \Gamma, x : A, y : B \vdash e : (x \sim_{1_A} a) \simeq (b \sim_{1_B} y) \end{array}}{\Gamma \vdash \text{univ}([x : A]b, [y : B]a, [x : A, y : B]e) : A \simeq B} \end{array}$$

Each universe is itself an object of the next universe; thus  $\mathbf{1} : \mathbf{Prop} : \mathbf{Set} : \mathbf{Groupoid}$ . We also have the following definitional equalities:  $\phi \sim_{1_{\mathbf{Prop}}} \psi \stackrel{\text{def}}{=} \phi \simeq \psi$ ,  $A \sim_{1_{\mathbf{Set}}} B \stackrel{\text{def}}{=} A \simeq B$ . As well as the normal operation of substitution, we have an operation of *path substitution*:

$$\frac{\Gamma, x : A \vdash b : B \quad \Gamma \vdash e : a \sim_{1_A} a'}{\Gamma \vdash b[x//e] : b[x/a] =_{B[x//e]} b[x/a']}$$

The system  $\lambda \simeq_2$  enjoys the following properties. Univalence holds definitionally; that is, an equality between types  $A \simeq B$  is exactly (definitionally) the type of equivalences between  $A$  and  $B$ . Also, transport respects reflexivity and composition definitionally.

This type theory has been formalised in Agda, using the method of the system **Kipling** from McBride [5]. The method ensures that, if  $s$  and  $t$  are definitionally equal expressions in  $\lambda \simeq_2$ , then  $\llbracket s \rrbracket$  and  $\llbracket t \rrbracket$  are definitionally equal objects in Agda. We interpret each context with a groupoid in Agda; that is, we define a type `data Cx : Set1` of contexts, and functions

```
[_]C : Cx → Set1
EQC : ∀ Γ → [ Γ ]C → [ Γ ]C → Set
EQC2 : ∀ {Γ} {a1 a2 b1 b2 : [ Γ ]C} →
  EQC Γ a1 a2 → EQC Γ b1 b2 → EQC Γ a1 b1 → EQC Γ a2 b2 → Set
```

The formalisation is available online at <https://github.com/radams78/Equality2>.

**Related Work** An earlier version of this system was presented in [2]. In this talk, we also give semantics to this system in Agda’s type theory extended with a native type of groupoids, and show how the syntax and semantics are formalised in Agda.

Our system is closely related to the system PHOML (Predicative Higher-Order Minimal Logic) presented in [1]. The system  $\lambda \simeq_2$  can be seen as an extension of PHOML with groupoids, and with a univalent equality for both sets and groupoids.

Cubical type theory [3, 4] has a very similar motivation to this work, and also offers a type theory with univalence and a computational interpretation. One difference with our system is that transport across the identity path is identity, and transport across  $p \bullet q$  is the composition of transport across  $p$  with transport across  $q$ , up to definitional equality. In cubical type theory, these equations only hold up to propositional equality.

## References

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