

Summary of Halmos' Naive Set Theory

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Chapter 1

Primitive Terms and Axioms

Let there be *sets*. We assume that everything is a set.

Let there be a binary relation of *membership*, \in . If $x \in A$ we say that x *belongs* to A , x is an *element* of A , or x is *contained* in A . If this does not hold we write $x \notin A$.

Definition 1.1 (Empty). A set is *empty* iff it has no elements; otherwise it is *nonempty*.

Definition 1.2 (Disjoint). Two sets A and B are *disjoint* iff there is no x such that $x \in A$ and $x \in B$.

Definition 1.3 (Pairwise Disjoint). We say the elements of a set x are *pairwise disjoint* iff, for all $y, z \in x$, either y and z are disjoint or $y = z$.

Definition 1.4 (Subset). Let A and B be sets. We say that A is a *subset* of B , or B *includes* A , and write $A \subseteq B$ or $B \supseteq A$, iff every element of A is an element of B .

Axiom 1.5 (Axiom of Extensionality). *If two sets have the same elements then they are equal.*

Axiom 1.6 (Axiom of Regularity). *For any nonempty set A , there exists $m \in A$ such that m and A are disjoint.*

Axiom 1.7 (Axiom of Union). *For every set A , there exists a set that contains all the elements that belong to at least one element of A .*

Axiom 1.8 (Axiom of Pairing). *For any two sets, there exists a set that they both belong to.*

Axiom Schema 1.9 (Axiom Schema of Replacement). *For any property $P(x, y)$, the following is an axiom:*

Let A be a set such that, for all $x \in A$, there exists at most one y such that $P(x, y)$. Then there exists a set B whose elements are exactly those sets y such that $\exists x \in A. P(x, y)$.

Axiom 1.10 (Axiom of Infinity). *There exists a set I that has an empty element and such that, for all $x \in I$, there exists $y \in I$ such that the elements of y are exactly x and the elements of x .*

Axiom 1.11 (Power Set Axiom). *For any set A , there exists a set B such that every subset of A belongs to B .*

Axiom 1.12 (Axiom of Choice). *Let A be a set of nonempty, pairwise disjoint sets. Then there exists a set C such that, for all $x \in A$, there is exactly one y such that $y \in C$ and $y \in x$.*

Chapter 2

Basic Properties and Operations on Sets

Axiom 2.1 (Axiom of Comprehension, Aussonderungsaxiom). *To every set A and to every condition $S(x)$ there corresponds a set B whose elements are exactly those elements x of A for which $S(x)$ holds.*

Definition 2.2. Given a set A and a condition $S(x)$, we write $\{x \in A : S(x)\}$ for the set whose elements are exactly those elements x of A for which $S(x)$ holds.

PROOF: This exists by the Axiom of Comprehension and is unique by the Axiom of Extensionality. \square

Definition 2.3 ((Unordered) Pair). For any sets a and b , the *(unordered) pair* $\{a, b\}$ is the set whose elements are just a and b .

PROOF: This exists by the Axioms of Pairing and Comprehension, and is unique by the Axiom of Extensionality. \square

Axiom 2.4 (Power Set Axiom). *For any set A , there exists a set that contains all the subsets of A .*

Definition 2.5 (Empty). A set is *empty* iff it has no elements; otherwise it is *non-empty*.

Axiom 2.6 (Axiom of Infinity). *There exists a set I such that:*

- *I has an element that is empty*
- *for all $x \in I$, there exists $y \in I$ such that the elements of y are exactly x and the elements of x .*

Definition 2.7 (Ordered Pair). For any sets a and b , the *ordered pair* (a, b) is defined by

$$(a, b) := \{\{a\}, \{a, b\}\} .$$

Definition 2.8 (Power Set). For any set A , the *power set* of A , $\mathcal{P}A$, is the set whose elements are exactly the subsets of A .

PROOF: This exists by the Power Set Axiom and Axiom of Comprehension, and is unique by the Axiom of Extensionality. \square

Definition 2.9 (Cartesian Product). For any sets A and B , the *Cartesian product* $A \times B$ is

$$A \times B := \{p \in \mathcal{P}\mathcal{P}(A \cup B) : \exists a \in A. \exists b \in B. p = (a, b)\} .$$

Definition 2.10 (Relation). A *relation* is a set of ordered pairs.

If R is a relation, we write xRy for $(x, y) \in R$.

Given sets X and Y , a relation *between X and Y* is a subset of $X \times Y$.

Given a set X , a relation *on X* is a relation between X and X .

Definition 2.11 (Function). Let X and Y be sets. A *function*, *map*, *mapping*, *transformation* or *operator f from X to Y* , $f : X \rightarrow Y$, is a relation f between X and Y such that, for all $x \in X$, there exists a unique $f(x) \in Y$, called the *value* of f at the *argument* x , such that $(x, f(x)) \in f$.

Definition 2.12 (Family). Let I and X be sets. A *family* of elements of X indexed by I is a function $a : I \rightarrow X$. We write a_i for $a(i)$, and $\{a_i\}_{i \in I}$ for a .

Definition 2.13 (Cartesian Product of a Family of Sets). Let $\{A_i\}_{i \in I}$ be a family of sets. The *Cartesian product* $\times_{i \in I} A_i$ is the set of all families $\{a_i\}_{i \in I}$ such that $\forall i \in I. a_i \in A_i$.

We write A^I for $\times_{i \in I} A$.

Axiom 2.14 (Axiom of Choice). *The Cartesian product of a non-empty family of non-empty sets is non-empty.*

Axiom 2.15 (Axiom of substitution). *If $S(a, b)$ is a sentence such that for each a in A the set $\{b : S(a, b)\}$ can be formed, then there exists a function F with domain A such that $F(a) = \{b : S(a, b)\}$ for each a in A .*

2.1 The Subset Relation

Theorem 2.16. *For any set A , we have $A \subseteq A$.*

PROOF: Every element of A is an element of A . \square

Theorem 2.17. *For any sets A , B and C , if $A \subseteq B$ and $B \subseteq C$ then $A \subseteq C$.*

PROOF: If every element of A is an element of B , and every element of B is an element of C , then every element of A is an element of C . \square

Theorem 2.18. *For any sets A and B , if $A \subseteq B$ and $B \subseteq A$ then $A = B$.*

PROOF: If every element of A is an element of B , and every element of B is an element of A , then A and B have the same elements, and therefore are equal by the Axiom of Extensionality. \square

Definition 2.19 (Proper Subset). Let A and B be sets. We say that A is a *proper* subset of B , or B *properly* includes A , and write $A \subsetneq B$ or $B \supsetneq A$, iff $A \subseteq B$ and $A \neq B$.

2.2 Comprehension Notation

Theorem 2.20. *There is no set that contains every set.*

PROOF:

$\langle 1 \rangle 1$. LET: A be a set.

PROVE: There exists a set B such that $B \notin A$.

$\langle 1 \rangle 2$. LET: $B = \{x \in A : x \notin x\}$

$\langle 1 \rangle 3$. If $B \in A$ then we have $B \in B$ if and only if $B \notin B$.

$\langle 1 \rangle 4$. $B \notin A$

\square

2.3 The Empty Set

Theorem 2.21. *There exists a set with no elements.*

PROOF: Immediate from the Axiom of Infinity. \square

Definition 2.22 (Empty Set). The *empty set* \emptyset is the set with no elements.

Theorem 2.23. *For any set A we have $\emptyset \subset A$.*

PROOF: Vacuous. \square

2.4 Unordered Pairs

Definition 2.24 (Singleton). For any set a , the *singleton* $\{a\}$ is defined to be $\{a, a\}$.

2.5 Unions

Definition 2.25 (Union). For any set \mathcal{C} , the *union* of \mathcal{C} , $\bigcup \mathcal{C}$, is the set whose elements are the elements of the elements of \mathcal{C} .

We write $\bigcup_{X \in \mathcal{A}} t[X]$ for $\bigcup \{t[X] \mid X \in \mathcal{A}\}$.

PROOF: This exists by the Union Axiom and Comprehension Axiom, and is unique by the Axiom of Extensionality. \square

Proposition 2.26.

$$\bigcup \emptyset = \emptyset$$

PROOF: There is no set that is an element of an element of \emptyset . \square

Proposition 2.27. *For any set A , we have $\bigcup\{A\} = A$.*

PROOF: For any x , we have x is an element of an element of $\{A\}$ if and only if x is an element of A . \square

Definition 2.28. We write $A \cup B$ for $\bigcup\{A, B\}$.

Proposition 2.29. *For any set A , we have $A \cup \emptyset = A$.*

PROOF: $x \in A \cup \emptyset$ iff $x \in A$ or $x \in \emptyset$, iff $x \in A$. \square

Proposition 2.30 (Idempotence). *For any set A , we have $A \cup A = A$.*

PROOF: $x \in A$ or $x \in A$ is equivalent to $x \in A$. \square

Proposition 2.31. *For any sets A and B , we have $A \subseteq B$ if and only if $A \cup B = B$.*

PROOF: For any x , the statement "if $x \in A$ then $x \in B$ " is equivalent to " $x \in A$ or $x \in B$ if and only if $x \in B$ ". \square

Proposition 2.32. *For any sets a and b , we have $\{a\} \cup \{b\} = \{a, b\}$.*

PROOF: Immediate from definitions. \square

2.6 Intersections

Definition 2.33 (Intersection). For any sets A and B , the *intersection* $A \cap B$ is defined to be $\{x \in A : x \in B\}$.

Proposition 2.34. *For any set A , we have $A \cap \emptyset = \emptyset$.*

PROOF: There is no x such that $x \in A$ and $x \in \emptyset$. \square

Proposition 2.35. *For any set A , we have*

$$A \cap A = A.$$

PROOF: We have $x \in A$ and $x \in A$ if and only if $x \in A$. \square

Proposition 2.36. *For any sets A and B , we have $A \subseteq B$ if and only if $A \cap B = A$.*

PROOF: For any x , the statement "if $x \in A$ then $x \in B$ " is equivalent to " $x \in A$ and $x \in B$ if and only if $x \in A$ ". \square

Proposition 2.37. *For any sets A , B and C , we have $C \subseteq A$ if and only if $(A \cap B) \cup C = A \cap (B \cup C)$.*

PROOF: The statement "if $x \in C$ then $x \in A$ " is equivalent to the statement " $((x \in A \wedge x \in B) \vee x \in C) \Leftrightarrow (x \in A \wedge (x \in B \vee x \in C))$ ". \square

Definition 2.38 (Disjoint). Two sets A and B are *disjoint* if and only if $A \cap B = \emptyset$.

Definition 2.39 (Pairwise Disjoint). Let A be a set. We say the elements of A are *pairwise disjoint* if and only if, for all $x, y \in A$, if $x \cap y \neq \emptyset$ then $x = y$.

Definition 2.40 (Intersection). For any nonempty set \mathcal{C} , the *intersection* of \mathcal{C} , $\bigcap \mathcal{C}$, is the set that contains exactly those sets that belong to every element of \mathcal{C} .

We write $\bigcap_{X \in \mathcal{A}} t[X]$ for $\bigcap \{t[X] \mid X \in \mathcal{A}\}$.

PROOF:

$\langle 1 \rangle 1$. LET: \mathcal{C} be a nonempty set.

$\langle 1 \rangle 2$. There exists a set I whose elements are exactly the sets that belong to every element of \mathcal{C} .

PROOF: Pick $A \in \mathcal{C}$, and take $I = \{x \in A : \forall X \in \mathcal{C}. x \in X\}$.

$\langle 1 \rangle 3$. For any sets I, J , if the elements of I and J are exactly the sets that belong to every element of \mathcal{C} then $I = J$.

PROOF: Axiom of Extensionality.

\square

2.7 Unordered Triples

Definition 2.41 ((Unordered) Triple). Given sets a_1, \dots, a_n , define the (*unordered*) *n-tuple* $\{a_1, \dots, a_n\}$ to be

$$\{a_1, \dots, a_n\} := \{a_1\} \cup \dots \cup \{a_n\} .$$

2.8 Relative Complements

Definition 2.42 (Relative Complement). For any sets A and B , the *difference* or *relative complement* $A - B$ is defined to be

$$A - B := \{x \in A : x \notin B\} .$$

Proposition 2.43. For any sets A and E , we have $A \subseteq E$ if and only if

$$E - (E - A) = A$$

PROOF:

$\langle 1 \rangle 1$. LET: A and E be sets.

$\langle 1 \rangle 2$. If $A \subseteq E$ then $E - (E - A) = A$

$\langle 2 \rangle 1$. ASSUME: $A \subseteq E$

$\langle 2 \rangle 2$. $E - (E - A) \subseteq A$

PROOF: If $x \in E$ and $x \notin E - A$ then $x \in A$.

$\langle 2 \rangle 3$. $A \subseteq E - (E - A)$

PROOF: If $x \in A$ then $x \in E$ and $x \notin E - A$.

$\langle 1 \rangle 3$. If $E - (E - A) = A$ then $A \subseteq E$.

PROOF: Since $E - (E - A) \subseteq E$.

□

Proposition 2.44. *For any set E we have*

$$E - \emptyset = E$$

PROOF: $x \in E$ if and only if $x \in E$ and $x \notin \emptyset$. □

Proposition 2.45. *For any set E we have*

$$E - E = \emptyset .$$

PROOF: There is no x such that $x \in E$ and $x \notin E$. □

Proposition 2.46. *For any sets A and E , we have*

$$A \cap (E - A) = \emptyset .$$

PROOF: There is no x such that $x \in A$ and $x \in E - A$. □

Proposition 2.47. *Let A and E be sets. Then $A \subseteq E$ if and only if*

$$A \cup (E - A) = E .$$

PROOF:

$\langle 1 \rangle 1$. LET: A and E be sets.

$\langle 1 \rangle 2$. If $A \subseteq E$ then $A \cup (E - A) = E$.

$\langle 2 \rangle 1$. ASSUME: $A \subseteq E$

$\langle 2 \rangle 2$. $A \cup (E - A) \subseteq E$

PROOF: If $x \in A$ or $x \in E - A$ then $x \in E$.

$\langle 2 \rangle 3$. $E \subseteq A \cup (E - A)$

PROOF: If $x \in E$ then either $x \in A$ or $x \notin A$. In the latter case, $x \in E - A$.

$\langle 1 \rangle 3$. If $A \cup (E - A) = E$ then $A \subseteq E$

PROOF: Since $A \subseteq A \cup (E - A)$.

□

Proposition 2.48. *Let A , B and E be sets. Then:*

1. *If $A \subseteq B$ then $E - B \subseteq E - A$.*

2. *If $A \subseteq E$ and $E - B \subseteq E - A$ then $A \subseteq B$.*

PROOF:

$\langle 1 \rangle 1$. LET: A , B and E be sets.

$\langle 1 \rangle 2$. If $A \subseteq B$ then $E - B \subseteq E - A$.

PROOF: If $A \subseteq B$, $x \in E$ and $x \notin B$, then we have $x \in E$ and $x \notin A$.

⟨1⟩3. If $A \subseteq E$ and $E - B \subseteq E - A$ then $A \subseteq B$.

⟨2⟩1. ASSUME: $A \subseteq E$

⟨2⟩2. ASSUME: $E - B \subseteq E - A$

⟨2⟩3. LET: $x \in A$

⟨2⟩4. $x \in E$

⟨2⟩5. $x \notin E - A$

⟨2⟩6. $x \notin E - B$

⟨2⟩7. $x \in B$

□

Example 2.49. We cannot remove the hypothesis $A \subseteq E$ in item 2 above. Let $E = \emptyset$, $A = \{\emptyset\}$ and $B = \emptyset$. Then $E - B = E - A = \emptyset$ but $A \not\subseteq B$.

Proposition 2.50 (De Morgan's Law). *For any sets A , B and E , we have $E - (A \cup B) = (E - A) \cap (E - B)$.*

PROOF: $(x \in E \wedge \neg(x \in A \vee x \in B)) \Leftrightarrow (x \in E \wedge x \notin A \wedge x \in E \wedge x \notin B)$. □

Proposition 2.51 (De Morgan's Law). *For any sets A , B and E , we have $E - (A \cap B) = (E - A) \cup (E - B)$.*

PROOF: $(x \in E \vee \neg(x \in A \wedge x \in B)) \Leftrightarrow (x \in E \wedge x \notin A) \vee (x \in E \wedge x \notin B)$. □

Proposition 2.52. *For any sets A , B and E , if $A \subseteq E$ then*

$$A - B = A \cap (E - B) .$$

PROOF: If $A \subseteq E$ then we have $(x \in A \wedge x \notin B) \Leftrightarrow (x \in A \wedge x \in E \wedge x \notin B)$. □

Proposition 2.53. *For any sets A and B , we have $A \subseteq B$ if and only if $A - B = \emptyset$.*

PROOF: Both are equivalent to the statement that there is no x such that $x \in A$ and $x \notin B$. □

Proposition 2.54. *For any sets A and B , we have*

$$A - (A - B) = A \cap B .$$

PROOF: $(x \in A \wedge \neg(x \in A \wedge x \notin B)) \Leftrightarrow x \in A \wedge x \in B$. □

Proposition 2.55. *For any sets A , B and C , we have*

$$A \cap (B - C) = (A \cap B) - (A \cap C) .$$

PROOF: $(x \in A \wedge x \in B \wedge x \notin C) \Leftrightarrow (x \in A \wedge x \in B \wedge \neg(x \in A \wedge x \in C))$. □

Proposition 2.56. *For any sets A , B , C and E , if $(A \cap B) - C \subseteq E$ then we have*

$$A \cap B \subseteq (A \cap C) \cup (B \cap (E - C)) .$$

PROOF:

$\langle 1 \rangle 1.$ LET: $x \in A \cap B$
 PROVE: $x \in (A \cap C) \cup (B \cap (E - C))$
 $\langle 1 \rangle 2.$ CASE: $x \in C$
 PROOF: Then $x \in A \cap C$.
 $\langle 1 \rangle 3.$ CASE: $x \notin C$
 PROOF: Then $x \in E$ and so $x \in B \cap (E - C)$.
 \square

Proposition 2.57. *For any sets A, B, C and E , we have*

$$(A \cup C) \cap (B \cup (E - C)) \subseteq A \cup B .$$

PROOF: The statement $(x \in A \vee x \in C) \wedge (x \in B \vee (x \in E \wedge x \notin C))$ implies $x \in A \vee x \in B$. \square

Proposition 2.58 (De Morgan's Law). *Let E be a set and \mathcal{C} a nonempty set. Then*

$$E - \bigcup_{X \in \mathcal{C}} X = \bigcap_{X \in \mathcal{C}} (E - X) .$$

PROOF: Easy. \square

Proposition 2.59 (De Morgan's Law). *Let E be a set and \mathcal{C} a nonempty set. Then*

$$E - \bigcap_{X \in \mathcal{C}} X = \bigcup_{X \in \mathcal{C}} (E - X) .$$

PROOF: Easy. \square

2.9 Symmetric Difference

Definition 2.60 (Symmetric Difference). For any sets A and B , the *symmetric difference* $A + B$ is defined to be

$$A + B := (A - B) \cup (B - A) .$$

Proposition 2.61. *For any sets A and B , we have*

$$A + B = B + A$$

PROOF: From the commutativity of union. \square

Proposition 2.62. *For any sets A, B and C , we have*

$$A + (B + C) = (A + B) + C .$$

PROOF: Each is the set of all x that belong to either exactly one or all three of A, B and C . \square

Proposition 2.63. *For any set A , we have*

$$A + \emptyset = A .$$

PROOF:

$$\begin{aligned} A + \emptyset &= (A - \emptyset) \cup (\emptyset - A) \\ &= A \cup \emptyset \\ &= A \end{aligned}$$

□

Proposition 2.64. *For any set A we have*

$$A + A = \emptyset .$$

PROOF:

$$\begin{aligned} A + A &= (A - A) \cup (A - A) \\ &= \emptyset \cup \emptyset \\ &= \emptyset \end{aligned}$$

□

2.10 Power Sets

Proposition 2.65.

$$\mathcal{P}\emptyset = \{\emptyset\}$$

PROOF: The only subset of \emptyset is \emptyset . □

Proposition 2.66. *For any set a , we have*

$$\mathcal{P}\{a\} = \{\emptyset, \{a\}\} .$$

PROOF: The only subsets of $\{a\}$ are \emptyset and $\{a\}$. □

Proposition 2.67. *For any sets a and b , we have*

$$\mathcal{P}\{a, b\} = \{\emptyset, \{a\}, \{b\}, \{a, b\}\} .$$

PROOF: The only subsets of $\{a, b\}$ are \emptyset , $\{a\}$, $\{b\}$ and $\{a, b\}$. □

Proposition 2.68. *For any nonempty set \mathcal{C} we have*

$$\bigcap_{X \in \mathcal{C}} \mathcal{P}X = \mathcal{P}\left(\bigcap \mathcal{C}\right) .$$

PROOF:

$$\begin{aligned} x \in \bigcup_{X \in \mathcal{C}} \mathcal{P}X &\Leftrightarrow \forall X \in \mathcal{C}. x \subseteq X \\ &\Leftrightarrow \forall X \in \mathcal{C}. \forall y \in x. y \in X \\ &\Leftrightarrow \forall y \in x. \forall X \in \mathcal{C}. y \in X \\ &\Leftrightarrow x \subseteq \bigcap \mathcal{C} \end{aligned}$$

□

Proposition 2.69. *For any set \mathcal{C} we have*

$$\bigcup_{X \in \mathcal{C}} \mathcal{P}X \subseteq \mathcal{P} \bigcup \mathcal{C} .$$

PROOF: If there exists $X \in \mathcal{C}$ such that $x \subseteq X$ then $x \subseteq \bigcup \mathcal{C}$. \square

Proposition 2.70. *For any set E , we have*

$$\bigcap \mathcal{P}E = \emptyset .$$

PROOF: Since $\emptyset \in \mathcal{P}E$. \square

Proposition 2.71. *For any sets E and F , if $E \subseteq F$ then $\mathcal{P}E \subseteq \mathcal{P}F$.*

PROOF: If $E \subseteq F$ and $X \subseteq E$ then $X \subseteq F$. \square

Chapter 3

Relations and Functions

3.1 Ordered Pairs

Proposition 3.1. *For any sets a, b, x and y , if $(a, b) = (x, y)$ then $a = x$ and $b = y$.*

PROOF:

$\langle 1 \rangle 1$. LET: a, b, x and y be sets.

$\langle 1 \rangle 2$. ASSUME: $(a, b) = (x, y)$

$\langle 1 \rangle 3$. $a = x$

PROOF: $\{a\} = \bigcap(a, b) = \bigcap(x, y) = \{x\}$.

$\langle 1 \rangle 4$. $\{a, b\} = \{x, y\}$

$\langle 1 \rangle 5$. CASE: $a = b$

$\langle 2 \rangle 1$. $x = y$

PROOF: Since $\{x, y\} = \{a, b\}$ is a singleton.

$\langle 2 \rangle 2$. $b = y$

PROOF: $b = a = x = y$

$\langle 1 \rangle 6$. CASE: $a \neq b$

$\langle 2 \rangle 1$. $x \neq y$

PROOF: Since $\{x, y\} = \{a, b\}$ is not a singleton.

$\langle 2 \rangle 2$. $b = y$

PROOF: $\{b\} = \{a, b\} - \{a\} = \{x, y\} - \{x\} = \{y\}$.

□

Proposition 3.2. *For any sets A, B and X , we have*

$$(A - B) \times X = (A \times X) - (B \times X) .$$

PROOF: Easy. □

Proposition 3.3. *For any sets A and B , we have $A \times B = \emptyset$ if and only if $A = \emptyset$ or $B = \emptyset$.*

PROOF: Easy. □

Proposition 3.4. For any sets A , B , X and Y , if $A \subseteq X$ and $B \subseteq Y$ then $A \times B \subseteq X \times Y$. The converse holds assuming $A \neq \emptyset$ and $B \neq \emptyset$.

PROOF: Easy. \square

3.2 Relations

Definition 3.5 (Domain). The *domain* of a relation R is the set

$$\text{dom } R := \left\{ x \in \bigcup \bigcup R : \exists y. (x, y) \in R \right\} .$$

Definition 3.6 (Range). The *range* of a relation R is the set

$$\text{ran } R := \left\{ y \in \bigcup \bigcup R : \exists x. (x, y) \in R \right\} .$$

Definition 3.7 (Reflexive). Let R be a relation on X . Then R is *reflexive* iff, for all $x \in X$, we have xRx .

Definition 3.8 (Symmetric). Let R be a relation on X . Then R is *symmetric* iff, whenever xRy , then yRx .

Definition 3.9 (Antisymmetric). A relation R is *antisymmetric* iff, whenever xRy and yRx , then $x = y$.

Definition 3.10 (Transitive). Let R be a relation on X . Then R is *transitive* iff, whenever xRy and yRz , then xRz .

Definition 3.11 (Identity Relation). For any set X , the *identity relation* I_X on X is

$$I_X = \{(x, x) : x \in X\} .$$

3.3 Composition

Definition 3.12 (Composition). Let R be a relation between X and Y , and S a relation between Y and Z . The *composite* or *relative product* $S \circ R = SR$ is the relation between X and Z defined by

$$x(S \circ R)z \Leftrightarrow \exists y \in Y (xRy \wedge ySz) .$$

Proposition 3.13. Let R be a relation between X and Y , S a relation between Y and Z , and T a relation between Z and W . Then

$$T(SR) = (TS)R .$$

PROOF: Easy. \square

Example 3.14. Composition of relations is not commutative in general. Let $X = \{a, b\}$ where $a \neq b$. Let $R = \{(a, a), (b, a)\}$ and $S = \{(a, b), (b, b)\}$. Then $SR = S$ but $RS = R \neq S$.

Proposition 3.15. A relation R is transitive if and only if $RR \subseteq R$.

PROOF: Easy. \square

3.4 Inverses

Definition 3.16 (Inverse). Let R be a relation between X and Y . The *inverse* or *converse* R^{-1} is the relation between Y and X defined by

$$yR^{-1}x \Leftrightarrow xRy .$$

Proposition 3.17. *For any relation R , we have*

$$\text{dom } R^{-1} = \text{ran } R .$$

PROOF: Easy. \square

Proposition 3.18. *For any relation R , we have*

$$\text{ran } R^{-1} = \text{dom } R .$$

PROOF: Easy. \square

Proposition 3.19. *Let R be a relation between X and Y , and S a relation between Y and Z . Then*

$$(SR)^{-1} = R^{-1}S^{-1} .$$

PROOF: Easy. \square

Proposition 3.20. *A relation R is symmetric if and only if $R \subseteq R^{-1}$.*

PROOF: Easy. \square

Proposition 3.21. *Let R be a relation between X and Y . Then*

$$I_Y R = R I_X = R .$$

PROOF: Easy. \square

Proposition 3.22. *A relation R on a set X is reflexive if and only if $I_X \subseteq R$.*

PROOF: Easy. \square

Proposition 3.23. *Let R be a relation on a set X . Then R is antisymmetric iff $R \cap R^{-1} \subseteq I_X$.*

PROOF: Easy. \square

3.5 Equivalence Relations

Definition 3.24 (Equivalence Relation). Let R be a relation on X . Then R is an *equivalence relation* iff it is reflexive, symmetric and transitive.

Definition 3.25 (Partition). Let X be a set. A *partition* of X is a pairwise disjoint set of nonempty subsets of X whose union is X .

Definition 3.26 (Equivalence Class). Let R be an equivalence relation on X . Let $x \in X$. The *equivalence class* of x with respect to R is

$$x/R := \{y \in X : xRy\} .$$

We write X/R for the set of all equivalence classes with respect to R .

Definition 3.27 (Induced). Let P be a partition of X . The relation *induced* by P is X/P where $x(X/P)y$ iff there exists $X \in P$ such that $x \in X$ and $y \in X$.

Theorem 3.28. Let R be an equivalence relation on X . Then X/R is a partition of X that induces the relation R .

PROOF: Easy. \square

Theorem 3.29. Let P be a partition of X . Then X/P is an equivalence relation on X , and $P = X/(X/P)$.

PROOF: Easy. \square

3.6 Functions

Definition 3.30 (Injective). A function $f : X \rightarrow Y$ is *one-to-one* or *injective* or an *injection* iff, for all $x, y \in X$, if $f(x) = f(y)$ then $x = y$. In this case, we write $f : X \rightarrowtail Y$.

Definition 3.31 (Surjective). Let $f : X \rightarrow Y$. We say f is *surjective*, or a *surjection*, or f maps X *onto* Y iff $\text{ran } f = Y$. In this case, we write $f : X \twoheadrightarrow Y$.

Definition 3.32 (Bijective). Let $f : X \rightarrow Y$. Then f is *bijective*, or a *bijection*, iff it is injective and surjective.

Definition 3.33 (Image). Let $f : X \rightarrow Y$ and $A \subseteq X$. The *image* of A under f is

$$f(A) := \{f(x) : x \in A\} .$$

Proposition 3.34. Let $f : X \rightarrow Y$ and $A \subseteq B \subseteq X$. Then $f(A) \subseteq f(B)$.

PROOF:

$\langle 1 \rangle 1$. LET: X and Y be sets.

$\langle 1 \rangle 2$. LET: $f : X \rightarrow Y$

$\langle 1 \rangle 3$. LET: $A, B \in \mathcal{P}X$ with $A \subseteq B$.

$\langle 1 \rangle 4$. LET: $y \in f(A)$

$\langle 1 \rangle 5$. PICK $x \in A$ such that $f(x) = y$

PROOF: $\langle 1 \rangle 4$

$\langle 1 \rangle 6$. $x \in B$

PROOF: $\langle 1 \rangle 3, \langle 1 \rangle 5$

$\langle 1 \rangle 7$. $y \in f(B)$

PROOF: $\langle 1 \rangle 5, \langle 1 \rangle 6$

□

Proposition 3.35. *Let $f : X \rightarrow Y$. Let $\mathcal{A} \subseteq \mathcal{P}X$. Then $f(\bigcup \mathcal{A}) = \bigcup_{A \in \mathcal{A}} f(A)$.*

PROOF:

$$\begin{aligned}
 y \in f\left(\bigcup \mathcal{A}\right) &\Leftrightarrow \exists x \in \bigcup \mathcal{A}. y = f(x) \\
 &\Leftrightarrow \exists x. \exists A \in \mathcal{A} (x \in A \wedge y = f(x)) \\
 &\Leftrightarrow \exists A \in \mathcal{A}. \exists x \in A. y = f(x) \\
 &\Leftrightarrow \exists A \in \mathcal{A}. y \in f(A) \\
 &\Leftrightarrow y \in \bigcup_{A \in \mathcal{A}} f(A) \quad \square
 \end{aligned}$$

Proposition 3.36. *Let $f : X \rightarrow Y$. Let \mathcal{A} be a nonempty subset of $\mathcal{P}X$. Then $f(\bigcap \mathcal{A}) \subseteq \bigcap_{A \in \mathcal{A}} f(A)$. Equality holds if f is injective.*

PROOF:

- ⟨1⟩1. LET: X and Y be sets.
- ⟨1⟩2. LET: $f : X \rightarrow Y$
- ⟨1⟩3. LET: \mathcal{A} be a nonempty subset of $\mathcal{P}X$.
- ⟨1⟩4. $f(\bigcap \mathcal{A}) \subseteq \bigcap_{A \in \mathcal{A}} f(A)$
 - ⟨2⟩1. LET: $y \in f(\bigcap \mathcal{A})$
 - ⟨2⟩2. PICK $x \in \bigcap \mathcal{A}$ such that $y = f(x)$
 - PROOF: ⟨2⟩1
 - ⟨2⟩3. LET: $A \in \mathcal{A}$
 - ⟨2⟩4. $x \in A$
 - PROOF: ⟨2⟩2, ⟨2⟩3
 - ⟨2⟩5. $y \in f(A)$
 - PROOF: ⟨2⟩2, ⟨2⟩4
- ⟨1⟩5. If f is injective then $f(\bigcap \mathcal{A}) = \bigcap_{A \in \mathcal{A}} f(A)$
 - ⟨2⟩1. ASSUME: f is injective.
 - ⟨2⟩2. LET: $y \in \bigcap_{A \in \mathcal{A}} f(A)$
 - ⟨2⟩3. PICK $A \in \mathcal{A}$
 - PROOF: \mathcal{A} is nonempty by ⟨1⟩3.
 - ⟨2⟩4. $y \in f(A)$
 - PROOF: ⟨2⟩2, ⟨2⟩3
 - ⟨2⟩5. PICK $x \in A$ such that $y = f(x)$
 - PROOF: ⟨2⟩4
 - ⟨2⟩6. $x \in \bigcap \mathcal{A}$
 - ⟨3⟩1. LET: $A' \in \mathcal{A}$
 - ⟨3⟩2. $y \in f(A')$
 - PROOF: ⟨2⟩2, ⟨3⟩1
 - ⟨3⟩3. PICK $x' \in A'$ such that $y = f(x')$
 - PROOF: ⟨3⟩2
 - ⟨3⟩4. $x = x'$
 - PROOF: ⟨2⟩1, ⟨2⟩5, ⟨3⟩3
 - ⟨3⟩5. $x \in A'$

PROOF: $\langle 3 \rangle 3, \langle 3 \rangle 4$
 $\langle 2 \rangle 7. y \in f(\bigcap \mathcal{A})$
 PROOF: $\langle 2 \rangle 5, \langle 2 \rangle 6$

□

Proposition 3.37. *Let X and Y be sets. Let $f : X \rightarrow Y$. Let $A, B \in \mathcal{P}X$. Then $f(A) - f(B) \subseteq f(A - B)$. Equality holds if f is injective.*

PROOF:

$\langle 1 \rangle 1$. LET: X and Y be sets.
 $\langle 1 \rangle 2$. LET: $f : X \rightarrow Y$
 $\langle 1 \rangle 3$. LET: $A, B \in \mathcal{P}X$
 $\langle 1 \rangle 4$. $f(A) - f(B) \subseteq f(A - B)$
 $\langle 2 \rangle 1$. LET: $y \in f(A) - f(B)$
 $\langle 2 \rangle 2$. $y \in f(A)$
 PROOF: $\langle 2 \rangle 1$
 $\langle 2 \rangle 3$. PICK $x \in A$ such that $y = f(x)$.
 PROOF: $\langle 2 \rangle 2$
 $\langle 2 \rangle 4$. $x \notin B$
 $\langle 3 \rangle 1$. ASSUME: for a contradiction $x \in B$
 $\langle 3 \rangle 2$. $y \in f(B)$
 PROOF: $\langle 2 \rangle 3, \langle 3 \rangle 1$
 $\langle 3 \rangle 3$. Q.E.D.
 PROOF: $\langle 2 \rangle 1$ and $\langle 3 \rangle 2$ form a contradiction.
 $\langle 2 \rangle 5$. $x \in A - B$
 PROOF: $\langle 2 \rangle 3, \langle 2 \rangle 4$
 $\langle 2 \rangle 6$. $y \in f(A - B)$
 PROOF: $\langle 2 \rangle 3, \langle 2 \rangle 5$
 $\langle 1 \rangle 5$. If f is injective then $f(A - B) = f(A) - f(B)$.
 $\langle 2 \rangle 1$. ASSUME: f is injective.
 $\langle 2 \rangle 2$. LET: $y \in f(A - B)$
 $\langle 2 \rangle 3$. PICK $x \in A - B$ such that $y = f(x)$
 PROOF: $\langle 2 \rangle 2$
 $\langle 2 \rangle 4$. $x \in A$
 PROOF: $\langle 2 \rangle 3$
 $\langle 2 \rangle 5$. $y \in f(A)$
 PROOF: $\langle 2 \rangle 3, \langle 2 \rangle 4$
 $\langle 2 \rangle 6$. $y \notin f(B)$
 $\langle 3 \rangle 1$. ASSUME: $y \in f(B)$
 $\langle 3 \rangle 2$. PICK $x' \in B$ such that $y = f(x')$
 PROOF: $\langle 3 \rangle 1$
 $\langle 3 \rangle 3$. $x = x'$
 PROOF: $\langle 2 \rangle 1, \langle 2 \rangle 3, \langle 3 \rangle 2$
 $\langle 3 \rangle 4$. $x \in B$
 PROOF: $\langle 3 \rangle 2, \langle 3 \rangle 3$
 $\langle 3 \rangle 5$. Q.E.D.
 PROOF: $\langle 2 \rangle 3$ and $\langle 3 \rangle 4$ form a contradiction.

$\langle 2 \rangle 7. y \in f(A) - f(B)$

PROOF: $\langle 2 \rangle 5, \langle 2 \rangle 6$

□

Definition 3.38 (Inclusion Map). Let Y be a set and $X \subseteq Y$. Then the *inclusion map* $i : X \hookrightarrow Y$ is the function defined by $i(x) = x$ for all $x \in X$.

Proposition 3.39. For any set X , the identity relation I_X is a function $X \rightarrow X$.

PROOF: Easy. □

Definition 3.40 (Restriction). Let $f : Y \rightarrow Z$ and $X \subseteq Y$. The *restriction* of f to X is the function $f \upharpoonright X : X \rightarrow Z$ defined by

$$(f \upharpoonright X)(x) = f(x) \quad (x \in X) .$$

Given sets X, Y and Z with $X \subseteq Y$, if $f : X \rightarrow Z$ and $g : Y \rightarrow Z$, we say g is an *extension* of f to Y iff $f = g \upharpoonright X$.

Definition 3.41 (Projection). Given sets X and Y , the *projection* maps $\pi_1 : X \times Y \rightarrow X$ and $\pi_2 : X \times Y \rightarrow Y$ are defined by

$$\pi_1(x, y) = x, \quad \pi_2(x, y) = y \quad (x \in X, y \in Y) .$$

Definition 3.42 (Canonical Map). Let X be a set and R an equivalence relation on X . The *canonical map* $\pi : X \rightarrow X/R$ is the map defined by $\pi(x) = x/R$.

Proposition 3.43. Let $f : X \rightarrow Y$. Then the following are equivalent:

1. f is one-to-one.
2. For all $A, B \subseteq X$, we have $f(A \cap B) = f(A) \cap f(B)$.
3. For all $A \subseteq X$, we have $f(X - A) \subseteq Y - f(A)$.

PROOF: Easy. □

Proposition 3.44. Let $f : X \rightarrow Y$. Then f maps X onto Y if and only if, for all $A \subseteq X$, we have $Y - f(A) \subseteq f(X - A)$.

PROOF: Easy. □

3.7 Families

Proposition 3.45 (Generalized Associative Law for Unions). Let $\{I_j\}_{j \in J}$ be a family of sets. Let $K = \bigcup_{j \in J} I_j$. Let $\{A_k\}_{k \in K}$ be a family of sets indexed by K . Then

$$\bigcup_{k \in K} A_k = \bigcup_{j \in J} \bigcup_{i \in I_j} A_i .$$

PROOF: Easy. \square

Proposition 3.46 (Generalized Commutative Law for Unions). *Let $\{I_j\}_{j \in J}$ be a family of sets. Let $f : J \rightarrow J$ be a one-to-one correspondence from J onto J . Then*

$$\bigcup_{j \in J} I_j = \bigcup_{j \in J} I_{f(j)} .$$

PROOF: Easy. \square

Proposition 3.47 (Generalized Associative Law for Intersections). *Let $\{I_j\}_{j \in J}$ be a nonempty family of nonempty sets. Let $K = \bigcup_{j \in J} I_j$. Let $\{A_k\}_{k \in K}$ be a family of sets indexed by K . Then*

$$\bigcap_{k \in K} A_k = \bigcap_{j \in J} \bigcap_{i \in I_j} A_i .$$

PROOF: Easy. \square

Proposition 3.48 (Generalized Commutative Law for Intersections). *Let $\{I_j\}_{j \in J}$ be a nonempty family of sets. Let $f : J \rightarrow J$ be a one-to-one correspondence from J onto J . Then*

$$\bigcap_{j \in J} I_j = \bigcap_{j \in J} I_{f(j)} .$$

PROOF: Easy. \square

Proposition 3.49. *Let B be a set and $\{A_i\}_{i \in I}$ a family of sets. Then*

$$B \cap \bigcup_{i \in I} A_i = \bigcup_{i \in I} (B \cap A_i)$$

PROOF: Easy. \square

Proposition 3.50. *Let B be a set and $\{A_i\}_{i \in I}$ a nonempty family of sets. Then*

$$B \cup \bigcap_{i \in I} A_i = \bigcap_{i \in I} (B \cup A_i)$$

PROOF: Easy. \square

Definition 3.51 (Projection). *Let $\{A_i\}_{i \in I}$ be a family of sets and $i \in I$. The projection function $\pi_i : \times_{i \in I} A_i \rightarrow A_i$ is defined by $\pi_i(a) = a_i$.*

Proposition 3.52. *Let $\{A_i\}_{i \in I}$ and $\{B_j\}_{j \in J}$ be families of sets. Then*

$$\left(\bigcup_{i \in I} A_i \right) \times \left(\bigcup_{j \in J} B_j \right) = \bigcup_{i \in I} \bigcup_{j \in J} (A_i \times B_j) .$$

PROOF: Easy. \square

Proposition 3.53. Let $\{A_i\}_{i \in I}$ and $\{B_j\}_{j \in J}$ be nonempty families of sets. Then

$$\left(\bigcap_{i \in I} A_i \right) \times \left(\bigcap_{j \in J} B_j \right) = \bigcap_{i \in I} \bigcap_{j \in J} (A_i \times B_j) .$$

PROOF: Easy. \square

Proposition 3.54. Let $f : X \rightarrow Y$. Let $\{A_i\}_{i \in I}$ be a family of subsets of X . Then

$$f \left(\bigcup_{i \in I} A_i \right) = \bigcup_{i \in I} f(A_i) .$$

PROOF: Easy. \square

Example 3.55. It is not true in general that, if $f : X \rightarrow Y$ and $\{A_i\}_{i \in I}$ is a nonempty family of subsets of X , then $f(\bigcap_{i \in I} A_i) = \bigcap_{i \in I} f(A_i)$.

Take $X = \{a, b\}$ and $Y = \{c\}$ where $a \neq b$. Take $I = \{i, j\}$ with $i \neq j$. Let $A_i = \{a\}$ and $A_j = \{b\}$. Let f be the unique function $X \rightarrow Y$. Then $f(\bigcap_{i \in I} A_i) = f(\emptyset) = \emptyset$ but $\bigcap_{i \in I} f(A_i) = \{c\}$.

3.8 Inverses and Composites of Functions

3.8.1 Inverse Image

Definition 3.56 (Inverse Image). Let $f : X \rightarrow Y$. Let B be a subset of Y . Then the *inverse image* of B under f is

$$f^{-1}(B) = \{x \in X : f(x) \in B\} .$$

Proposition 3.57. Let $f : X \rightarrow Y$. Let $B \subseteq Y$. Then

$$f(f^{-1}(B)) \subseteq B .$$

Equality holds if f is surjective.

PROOF:

- $\langle 1 \rangle 1$. LET: X and Y be sets.
- $\langle 1 \rangle 2$. LET: $f : X \rightarrow Y$
- $\langle 1 \rangle 3$. LET: $B \subseteq Y$
- $\langle 1 \rangle 4$. $f(f^{-1}(B)) \subseteq B$
 - $\langle 2 \rangle 1$. LET: $y \in f(f^{-1}(B))$
 - $\langle 2 \rangle 2$. PICK $x \in f^{-1}(B)$ such that $f(x) = y$
 - $\langle 2 \rangle 3$. $f(x) \in B$
 - $\langle 2 \rangle 4$. $y \in B$
- $\langle 1 \rangle 5$. If f is surjective then $f(f^{-1}(B)) = B$
 - $\langle 2 \rangle 1$. ASSUME: f is surjective.
 - $\langle 2 \rangle 2$. LET: $y \in B$

$\langle 2 \rangle 3$. PICK $x \in X$ such that $f(x) = y$

PROOF: $\langle 2 \rangle 1, \langle 2 \rangle 2$

$\langle 2 \rangle 4$. $x \in f^{-1}(B)$

PROOF: $\langle 2 \rangle 2, \langle 2 \rangle 3$

$\langle 2 \rangle 5$. $y \in f(f^{-1}(B))$

PROOF: $\langle 2 \rangle 3, \langle 2 \rangle 4$

□

Proposition 3.58. *Let $f : X \rightarrow Y$. Let $A \subseteq X$. Then*

$$A \subseteq f^{-1}(f(A)) .$$

Equality holds if f is one-to-one.

PROOF:

$\langle 1 \rangle 1$. LET: X and Y be sets.

$\langle 1 \rangle 2$. LET: $f : X \rightarrow Y$

$\langle 1 \rangle 3$. LET: $A \subseteq X$

$\langle 1 \rangle 4$. $A \subseteq f^{-1}(f(A))$

$\langle 2 \rangle 1$. LET: $x \in A$

$\langle 2 \rangle 2$. $f(x) \in f(A)$

$\langle 2 \rangle 3$. $x \in f^{-1}(f(x))$

$\langle 1 \rangle 5$. If f is injective then $f^{-1}(f(A)) = A$

$\langle 2 \rangle 1$. ASSUME: f is injective.

$\langle 2 \rangle 2$. LET: $x \in f^{-1}(f(A))$

$\langle 2 \rangle 3$. $f(x) \in f(A)$

PROOF: $\langle 2 \rangle 2$

$\langle 2 \rangle 4$. PICK $x' \in A$ such that $f(x') = f(x)$

PROOF: $\langle 2 \rangle 3$

$\langle 2 \rangle 5$. $x' = x$

PROOF: $\langle 2 \rangle 1, \langle 2 \rangle 4$

$\langle 2 \rangle 6$. $x \in A$

PROOF: $\langle 2 \rangle 4, \langle 2 \rangle 5$

□

Proposition 3.59. *Let $f : X \rightarrow Y$. Let $A \subseteq B \subseteq Y$. Then $f^{-1}(A) \subseteq f^{-1}(B)$.*

PROOF:

$\langle 1 \rangle 1$. LET: X and Y be sets.

$\langle 1 \rangle 2$. LET: $f : X \rightarrow Y$

$\langle 1 \rangle 3$. LET: $A \subseteq B \subseteq Y$

$\langle 1 \rangle 4$. LET: $x \in f^{-1}(A)$

$\langle 1 \rangle 5$. $f(x) \in A$

$\langle 1 \rangle 6$. $f(x) \in B$

$\langle 1 \rangle 7$. $x \in f^{-1}(B)$

□

Proposition 3.60. *Let $f : X \rightarrow Y$. Let $\mathcal{B} \subseteq Y$. Then*

$$f^{-1}\left(\bigcup \mathcal{B}\right) = \bigcup_{B \in \mathcal{B}} f^{-1}(B) .$$

PROOF:

$$\begin{aligned}
x \in f^{-1}\left(\bigcup \mathcal{B}\right) &\Leftrightarrow f(x) \in \bigcup \mathcal{B} \\
&\Leftrightarrow \exists B \in \mathcal{B}. f(x) \in B \\
&\Leftrightarrow \exists B \in \mathcal{B}. x \in f^{-1}(B) \\
&\Leftrightarrow x \in \bigcup_{B \in \mathcal{B}} f^{-1}(B) \quad \square
\end{aligned}$$

Proposition 3.61. *Let $f : X \rightarrow Y$. Let $\{B_i\}_{i \in I}$ be a nonempty family of subsets of Y . Then*

$$f^{-1}\left(\bigcap_{i \in I} B_i\right) = \bigcap_{i \in I} f^{-1}(B_i) .$$

PROOF: Easy. \square

Proposition 3.62. *Let $f : X \rightarrow Y$ and $B \subseteq Y$. Then $f^{-1}(Y - B) = X - f^{-1}(B)$.*

PROOF: Easy. \square

3.8.2 Inverse of a Function

Proposition 3.63. *Let $f : X \approx Y$. Then f^{-1} is a function, and is a bijection $f^{-1} : Y \approx X$.*

PROOF:

$\langle 1 \rangle 1$. LET: X and Y be sets.

$\langle 1 \rangle 2$. LET: $f : X \approx Y$

$\langle 1 \rangle 3$. f^{-1} is a function.

$\langle 2 \rangle 1$. LET: $(x, y), (x, z) \in f^{-1}$

$\langle 2 \rangle 2$. $(y, x), (z, x) \in f$

$\langle 2 \rangle 3$. $y = z$

PROOF: f is injective.

$\langle 1 \rangle 4$. $\text{dom } f^{-1} = Y$

PROOF: Proposition 3.17, $\langle 1 \rangle 2$

$\langle 1 \rangle 5$. $\text{ran } f^{-1} = X$

PROOF:

$$\begin{aligned}
x \in \text{ran } f^{-1} &\Leftrightarrow \exists y. (y, x) \in f^{-1} \\
&\Leftrightarrow \exists y. (x, y) \in f \\
&\Leftrightarrow x \in \text{dom } f \\
&\Leftrightarrow x \in X
\end{aligned}$$

$\langle 1 \rangle 6$. f^{-1} is injective.

$\langle 2 \rangle 1$. LET: $y, y' \in Y$

$\langle 2 \rangle 2$. ASSUME: $f^{-1}(y) = f^{-1}(y')$

$\langle 2 \rangle 3$. $y = y'$

PROOF: $y = f(f^{-1}(y)) = f(f^{-1}(y')) = y'$.

□

Proposition 3.64. *Let $f : X \rightarrow Y$ and $g : Y \rightarrow Z$. Then $gf : X \rightarrow Z$ and, for all $x \in X$, we have*

$$(g \circ f)(x) = g(f(x)) .$$

PROOF: Easy. □

Example 3.65. Example 3.14 shows that function composition is not commutative in general.

Proposition 3.66. *The composite of two injective functions is injective.*

PROOF:

⟨1⟩1. LET: $f : X \rightarrow Y$ and $g : Y \rightarrow Z$

⟨1⟩2. LET: $x, y \in X$

⟨1⟩3. ASSUME: $(g \circ f)(x) = (g \circ f)(y)$

⟨1⟩4. $g(f(x)) = g(f(y))$

⟨1⟩5. $f(x) = f(y)$

PROOF: g is injective.

⟨1⟩6. $x = y$

PROOF: f is injective.

□

Proposition 3.67. *Let $f : X \rightarrow Y$ and $g : Y \rightarrow Z$. If $g \circ f$ is injective then f is injective.*

PROOF: If $f(x) = f(y)$ then $g(f(x)) = g(f(y))$ and so $x = y$. □

Proposition 3.68. *The composite of two surjective functions is surjective.*

PROOF:

⟨1⟩1. LET: $f : X \rightarrow Y$ and $g : Y \rightarrow Z$

⟨1⟩2. LET: $z \in Z$

⟨1⟩3. PICK $y \in Y$ such that $g(y) = z$

PROOF: Since g is surjective.

⟨1⟩4. PICK $x \in X$ such that $f(x) = y$

PROOF: Since f is surjective.

⟨1⟩5. $(g \circ f)(x) = z$

□

Proposition 3.69. *Let $f : X \rightarrow Y$ and $g : Y \rightarrow Z$. If $g \circ f$ is surjective then g is surjective.*

PROOF: Let $z \in Z$. Pick $x \in X$ such that $g(f(x)) = z$. Then there exists y such that $g(y) = z$, namely $y = f(x)$. □

Proposition 3.70. *The composite of two bijective functions is bijective.*

PROOF: Propositions 3.66 and 3.68. □

Proposition 3.71. *Let X, Y and Z be sets. Let $f : X \rightarrow Y$ and $g : Y \rightarrow Z$. Let $A \subseteq Z$. Then*

$$(g \circ f)^{-1}(A) = f^{-1}(g^{-1}(A)) .$$

PROOF:

$\langle 1 \rangle 1$. LET: X, Y and Z be sets.

$\langle 1 \rangle 2$. LET: $f : X \rightarrow Y$ and $g : Y \rightarrow Z$

$\langle 1 \rangle 3$. LET: $A \subseteq Z$

$\langle 1 \rangle 4$. $(g \circ f)^{-1}(A) = f^{-1}(g^{-1}(A))$

PROOF:

$\langle 2 \rangle 1$. LET: $x \in X$

$\langle 2 \rangle 2$. $x \in (g \circ f)^{-1}(A) \Leftrightarrow x \in f^{-1}(g^{-1}(A))$

PROOF:

$$x \in (g \circ f)^{-1}(A) \Leftrightarrow (g \circ f)(x) \in A$$

$$\Leftrightarrow g(f(x)) \in A \quad (\text{Proposition 3.64, } \langle 1 \rangle 2, \langle 2 \rangle 1)$$

$$\Leftrightarrow f(x) \in g^{-1}(A)$$

$$\Leftrightarrow x \in f^{-1}(g^{-1}(A))$$

□

Proposition 3.72. *Let $f : X \approx Y$ and $g : Y \approx Z$. Then*

$$(gf)^{-1} = f^{-1}g^{-1} : Z \rightarrow X .$$

PROOF: Easy. □

Definition 3.73 (Left Inverse, Right Inverse). Let $f : X \rightarrow Y$ and $g : Y \rightarrow X$. Then g is a *left inverse* of f , and f is a *right inverse* of g , iff $g \circ f = I_X$.

Proposition 3.74. *Let $f : X \rightarrow Y$ and $g : Y \rightarrow X$. If $gf = I_X$ then f is one-to-one and g maps Y onto X .*

PROOF: Easy. □

Lemma 3.75. *Let $f : A \rightarrow B$. If there are functions $g : B \rightarrow A$ and $h : B \rightarrow A$ such that $\forall a \in A. g(f(a)) = a$ and $\forall b \in B. f(h(b)) = b$, then f is bijective and $g = h = f^{-1}$.*

PROOF:

$\langle 1 \rangle 1$. LET: A and B be sets.

$\langle 1 \rangle 2$. LET: $f : A \rightarrow B$ and $g, h : B \rightarrow A$

$\langle 1 \rangle 3$. ASSUME: $\forall a \in A. g(f(a)) = a$

$\langle 1 \rangle 4$. ASSUME: $\forall b \in B. f(h(b)) = b$

$\langle 1 \rangle 5$. f is injective.

PROOF: Proposition 3.74, $\langle 1 \rangle 2, \langle 1 \rangle 3$.

$\langle 1 \rangle 6$. f is surjective.

PROOF: Proposition 3.74, $\langle 1 \rangle 2, \langle 1 \rangle 4$.

$\langle 1 \rangle 7$. $g = h$

$\langle 2 \rangle 1$. LET: $b \in B$
 $\langle 2 \rangle 2$. $g(b) = h(b)$
 PROOF:

$$\begin{aligned} g(b) &= g(f(h(b))) && (\langle 1 \rangle 4, \langle 2 \rangle 1) \\ &= h(b) && (\langle 1 \rangle 3, \langle 1 \rangle 2, \langle 2 \rangle 1) \end{aligned}$$
 $\langle 1 \rangle 8$. $h = f^{-1}$
 $\langle 2 \rangle 1$. LET: $b \in B$
 $\langle 2 \rangle 2$. $f(h(b)) = b$
 PROOF: $\langle 1 \rangle 4, \langle 2 \rangle 1$
 $\langle 2 \rangle 3$. $h(b) = f^{-1}(b)$
 \square

3.9 Choice Functions

Definition 3.76 (Choice Function). A *choice function* for a set X is a function $f : \mathcal{P}X - \{\emptyset\} \rightarrow X$ such that $f(S) \in S$ for all S .

Proposition 3.77. *Every set has a choice function.*

PROOF: Given a nonempty set X , apply the Axiom of Choice to the family $\{S\}_{S \in \mathcal{P}X - \{\emptyset\}}$. \square

Proposition 3.78. *For any relation R , there exists a function $f \subseteq R$ such that $\text{dom } f = \text{dom } R$.*

PROOF:
 $\langle 1 \rangle 1$. LET: R be a relation.
 $\langle 1 \rangle 2$. PICK a choice function g for $\text{ran } R$.
 $\langle 1 \rangle 3$. LET: $f : \text{dom } R \rightarrow \text{ran } R$ be the function $f(x) = g(\{y \in \text{ran } R : xRy\})$
 $\langle 1 \rangle 4$. $f \subseteq R$ and $\text{dom } f = \text{dom } R$.
 \square

Proposition 3.79. *If \mathcal{C} is a set of pairwise disjoint nonempty sets, then there exists a set A such that, for all $C \in \mathcal{C}$, we have $A \cap C$ is a singleton.*

PROOF:
 $\langle 1 \rangle 1$. LET: f be a choice function for $\bigcup \mathcal{C}$
 $\langle 1 \rangle 2$. LET: $A = \{f(C) : C \in \mathcal{C}\}$
 $\langle 1 \rangle 3$. For all $C \in \mathcal{C}$ we have $A \cap C = \{f(C)\}$
 \square

Chapter 4

Equivalence

Definition 4.1 (Equivalent). Sets E and F are *equivalent*, $E \sim F$, iff there exists a one-to-one correspondence between them.

Proposition 4.2. *For any set X , equivalence is an equivalence relation on $\mathcal{P}X$.*

PROOF: Easy.

Theorem 4.3 (Schröder-Bernstein). *Let X and Y be sets. If there exist injective functions $X \rightarrow Y$ and $Y \rightarrow X$, then $X \sim Y$.*

PROOF:

- $\langle 1 \rangle 1$. LET: $f : X \rightarrow Y$ and $g : Y \rightarrow X$ be one-to-one.
- $\langle 1 \rangle 2$. ASSUME: w.l.o.g. $X \cap Y = \emptyset$
- $\langle 1 \rangle 3$. For $x \in X$, let us say that x is the *parent* of $f(x)$; and for $y \in Y$, let us say that y is the *parent* of $g(y)$.
- $\langle 1 \rangle 4$. For $z \in X \cup Y$, let the set of *descendants* of z be the intersection of all the subsets S of $X \cup Y$ such that $z \in S$ and, if $t \in S$ and t is the parent of u then $u \in S$.
- $\langle 1 \rangle 5$. LET: X_X be the set of all elements of X that are descendants of the elements of X that have no parent.
- $\langle 1 \rangle 6$. LET: X_Y be the set of all elements of X that are descendants of the elements of Y that have no parent.
- $\langle 1 \rangle 7$. LET: $X_\infty = X - X_X - X_Y$
- $\langle 1 \rangle 8$. LET: Y_X be the set of all elements of Y that are descendants of the elements of X that have no parent.
- $\langle 1 \rangle 9$. LET: Y_Y be the set of all elements of Y that are descendants of the elements of Y that have no parent.
- $\langle 1 \rangle 10$. LET: $Y_\infty = Y - Y_X - Y_Y$
- $\langle 1 \rangle 11$. $f|X_X : X_X \sim Y_X$
- $\langle 1 \rangle 12$. $g|Y_Y : Y_Y \sim X_Y$
- $\langle 1 \rangle 13$. $f|X_\infty : X_\infty \sim Y_\infty$
- $\langle 1 \rangle 14$. Define $h : X \rightarrow Y$ by $h(x) = g^{-1}(x)$ if $x \in X_Y$, and $f(x)$ if not.

15. $h : X \sim Y$
 \square

Theorem 4.4 (Cantor). *For any set X we have $X \not\sim \mathcal{P}X$.*

PROOF: If $f : X \rightarrow \mathcal{P}X$ then $\{x \in X : x \notin f(x)\}$ is a subset of X not in $\text{ran } f$. \square

Chapter 5

Order

5.1 Partial Orders

Definition 5.1 (Partial Order). A *partial order* on a set X is a relation on X that is reflexive, antisymmetric and transitive.

A *partially ordered set* or *poset* is a pair (X, \leq) such that \leq is a partial order on X . We write X for the poset (X, \leq) .

Given a partial order \leq , we write \geq for the inverse of \leq .

We write $x < y$ or $y > x$ for $x \leq y \wedge x \neq y$. When this holds, we say x is *less than* y or *smaller than* y ; and y is *greater than* x , *larger than* x .

Proposition 5.2. *For any set X , the relation \subseteq is a partial order on $\mathcal{P}X$.*

PROOF: From Theorems 2.16, 2.17 and 2.18. \square

Proposition 5.3. *Let X be a set. Let $<$ be a relation on X . Then there exists a partial order \leq on X such that*

$$\forall x, y \in X (x < y \Leftrightarrow x \leq y \wedge x \neq y)$$

if and only if $<$ is irreflexive and transitive. In this case, \leq is unique and is defined by

$$\forall x, y \in X (x \leq y \Leftrightarrow x < y \vee x = y) .$$

PROOF:

$\langle 1 \rangle 1$. LET: X be a set.

$\langle 1 \rangle 2$. LET: $<$ be a relation on X .

$\langle 1 \rangle 3$. If there exists a partial order \leq on X such that $\forall x, y \in X (x < y \Leftrightarrow x \leq y \wedge x \neq y)$, then $<$ is irreflexive.

PROOF: Trivial.

$\langle 1 \rangle 4$. If there exists a partial order \leq on X such that $\forall x, y \in X (x < y \Leftrightarrow x \leq y \wedge x \neq y)$, then $<$ is transitive.

$\langle 2 \rangle 1$. LET: \leq be a partial order on X .

$\langle 2 \rangle 2$. ASSUME: $\forall x, y \in X (x < y \Leftrightarrow x \leq y \wedge x \neq y)$

$\langle 2 \rangle 3$. LET: $x < y$ and $y < z$
 $\langle 2 \rangle 4$. $x \leq y$
PROOF: $\langle 2 \rangle 3$
 $\langle 2 \rangle 5$. $x \neq y$
PROOF: $\langle 2 \rangle 3$
 $\langle 2 \rangle 6$. $y \leq z$
PROOF: $\langle 2 \rangle 3$
 $\langle 2 \rangle 7$. $y \neq z$
PROOF: $\langle 2 \rangle 3$
 $\langle 2 \rangle 8$. $x \leq z$
PROOF: From $\langle 2 \rangle 4$ and $\langle 2 \rangle 6$ since \leq is transitive by $\langle 2 \rangle 1$.
 $\langle 2 \rangle 9$. $x \neq z$
 $\langle 3 \rangle 1$. ASSUME: for a contradiction $x = z$
 $\langle 3 \rangle 2$. $y \leq x$
PROOF: $\langle 2 \rangle 6$, $\langle 3 \rangle 1$
 $\langle 3 \rangle 3$. $x = y$
PROOF: From $\langle 2 \rangle 4$ and $\langle 3 \rangle 2$ since \leq is antisymmetric by $\langle 2 \rangle 1$.
 $\langle 3 \rangle 4$. Q.E.D.
PROOF: $\langle 2 \rangle 5$ and $\langle 3 \rangle 3$ form a contradiction.
 $\langle 2 \rangle 10$. $x < z$
PROOF: $\langle 2 \rangle 8$, $\langle 2 \rangle 9$
 $\langle 1 \rangle 5$. If there exists a partial order \leq on X such that $\forall x, y \in X (x < y \Leftrightarrow x \leq y \wedge x \neq y)$, then $\forall x, y \in X (x \leq y \Leftrightarrow x < y \vee x = y)$.
PROOF: Trivial.
 $\langle 1 \rangle 6$. If $<$ is irreflexive and transitive, then the relation \leq defined by $\forall x, y \in X (x \leq y \Leftrightarrow x < y \vee x = y)$ is a partial order on X .
 $\langle 2 \rangle 1$. ASSUME: $<$ is irreflexive.
 $\langle 2 \rangle 2$. ASSUME: $<$ is transitive.
 $\langle 2 \rangle 3$. LET: \leq be the relation defined by $\forall x, y \in X (x \leq y \Leftrightarrow x < y \vee x = y)$
 $\langle 2 \rangle 4$. \leq is reflexive.
PROOF: Immediate from $\langle 2 \rangle 3$.
 $\langle 2 \rangle 5$. \leq is transitive.
 $\langle 3 \rangle 1$. LET: $x, y, z \in X$
 $\langle 3 \rangle 2$. ASSUME: $x \leq y$ and $y \leq z$
 $\langle 3 \rangle 3$. $x < y$ or $x = y$
PROOF: $\langle 2 \rangle 4$, $\langle 3 \rangle 2$
 $\langle 3 \rangle 4$. $y < z$ or $y = z$
PROOF: $\langle 2 \rangle 4$, $\langle 3 \rangle 2$
 $\langle 3 \rangle 5$. CASE: $x < y$ and $y < z$
 $\langle 4 \rangle 1$. $x < z$
PROOF: $\langle 2 \rangle 2$
 $\langle 4 \rangle 2$. $x \leq z$
PROOF: $\langle 2 \rangle 4$, $\langle 4 \rangle 1$
 $\langle 3 \rangle 6$. CASE: $x = y$
PROOF: Then $x \leq z$ from $\langle 3 \rangle 2$
 $\langle 3 \rangle 7$. CASE: $y = z$

PROOF: Then $x \leq z$ from $\langle 3 \rangle 2$
 $\langle 2 \rangle 6$. \leq is antisymmetric.
 $\langle 3 \rangle 1$. LET: $x, y \in X$
 $\langle 3 \rangle 2$. ASSUME: $x \leq y$
 $\langle 3 \rangle 3$. ASSUME: $y \leq x$
 $\langle 3 \rangle 4$. $x < y$ or $x = y$
PROOF: $\langle 2 \rangle 4, \langle 3 \rangle 2$
 $\langle 3 \rangle 5$. $y < x$ or $y = x$
PROOF: $\langle 2 \rangle 4, \langle 3 \rangle 2$
 $\langle 3 \rangle 6$. $\neg(x < y \wedge y < x)$
 $\langle 4 \rangle 1$. ASSUME: for a contradiction $x < y$ and $y < x$
 $\langle 4 \rangle 2$. $x < x$
PROOF: $\langle 2 \rangle 2, \langle 4 \rangle 1$
 $\langle 4 \rangle 3$. Q.E.D.
PROOF: $\langle 2 \rangle 1$ and $\langle 4 \rangle 2$ form a contradiction.
 $\langle 3 \rangle 7$. $x = y$
PROOF: $\langle 3 \rangle 4, \langle 3 \rangle 5, \langle 3 \rangle 6$
 $\langle 1 \rangle 7$. If $<$ is irreflexive, then the relation \leq defined by $\forall x, y \in X (x \leq y \Leftrightarrow x < y \vee x = y)$ satisfies $\forall x, y \in X (x < y \Leftrightarrow x \leq y \wedge x \neq y)$.
PROOF: Trivial.
 \square

Proposition 5.4. *In a poset, we never have $x < y$ and $y < x$.*

PROOF: We would then have $x < x$ by transitivity, contradicting the irreflexivity of $<$. \square

Definition 5.5 ((Strict) Initial Segment). Let X be a poset and $a \in X$. The (strict) initial segment determined by a is

$$s(a) := \{x \in X : x < a\} .$$

Definition 5.6 (Weak Initial Segment). Let X be a poset and $a \in X$. The weak initial segment determined by a is

$$\bar{s}(a) := \{x \in X : x \leq a\} .$$

Definition 5.7 (Immediate Successor). Let X be a poset and $x, y \in X$. Then y is the (immediate) successor of x , and x is the (immediate) predecessor of y , iff $x < y$ and there is no z such that $x < z < y$.

Definition 5.8 (Least). Let X be a partial order and $a \in X$. Then a is least in X iff $\forall x \in X. a \leq x$.

Proposition 5.9. *A poset has at most one least element.*

PROOF: If a and b are least then $a \leq b$ and $b \leq a$, hence $a = b$. \square

Definition 5.10 (Greatest). Let X be a partial order and $a \in X$. Then a is greatest in X iff $\forall x \in X. x \leq a$.

Proposition 5.11. *A poset has at most one greatest element.*

PROOF: If a and b are greatest then $a \leq b$ and $b \leq a$, hence $a = b$. \square

Definition 5.12 (Minimal). Let X be a poset and $a \in X$. Then a is *minimal* iff there is no $x \in X$ such that $x < a$.

Definition 5.13 (Maximal). Let X be a poset and $a \in X$. Then a is *maximal* iff there is no $x \in X$ such that $a < x$.

Definition 5.14 (Lower Bound). Let X be a poset. Let $E \subseteq X$ and $a \in X$. Then a is a *lower bound* for E iff $\forall x \in E. a \leq x$.

Definition 5.15 (Upper Bound). Let X be a poset. Let $E \subseteq X$ and $a \in X$. Then a is an *upper bound* for E iff $\forall x \in E. x \leq a$.

Definition 5.16 (Greatest Lower Bound, Infimum). Let X be a poset. Let $E \subseteq X$ and $a \in X$. Then a is the *greatest lower bound* or *infimum* for E iff a is the greatest element in the set of lower bounds for E .

Definition 5.17 (Least Upper Bound, Supremum). Let X be a poset. Let $E \subseteq X$ and $a \in X$. Then a is the *least upper bound* or *supremum* for E iff a is the least element in the set of upper bounds for E .

5.2 Linear Orders

Definition 5.18 (Linear Order). A partial order \leq on a set X is a *linear order*, *total order* or *simple order* iff, for all $x, y \in X$, either $x \leq y$ or $y \leq x$. We then call the poset (X, \leq) a *linearly ordered set* or a *chain*.

Proposition 5.19. *Let R be a partial order on X . Then R is total if and only if $X^2 \subseteq R \cup R^{-1}$.*

PROOF: Easy. \square

Proposition 5.20. *Let X be a set and $<$ a relation on X . Then there exists a linear order \leq on X such that*

$$\forall x, y \in X. (x < y \Leftrightarrow x \leq y \wedge x \neq y)$$

if and only if $<$ is irreflexive, transitive, and:

$$\forall x, y \in X. (x < y \vee x = y \vee y < x) .$$

In this case, \leq is unique and is defined by

$$\forall x, y \in X. (x \leq y \Leftrightarrow x < y \vee x = y)$$

PROOF:

$\langle 1 \rangle$ 1. LET: X be a set.

- $\langle 1 \rangle 2$. LET: $<$ be a relation on X
 $\langle 1 \rangle 3$. If there exists a linear order \leq on X such that $\forall x, y \in X (x < y \Leftrightarrow x \leq y \wedge x \neq y)$, then $<$ is irreflexive.
 PROOF: Trivial.
 $\langle 1 \rangle 4$. If there exists a linear order \leq on X such that $\forall x, y \in X (x < y \Leftrightarrow x \leq y \wedge x \neq y)$, then $<$ is transitive.
 $\langle 2 \rangle 1$. LET: \leq be a linear order on X .
 $\langle 2 \rangle 2$. ASSUME: $\forall x, y \in X (x < y \Leftrightarrow x \leq y \wedge x \neq y)$
 $\langle 2 \rangle 3$. LET: $x < y$ and $y < z$
 $\langle 2 \rangle 4$. $x \leq y$
 PROOF: $\langle 2 \rangle 3$
 $\langle 2 \rangle 5$. $x \neq y$
 PROOF: $\langle 2 \rangle 3$
 $\langle 2 \rangle 6$. $y \leq z$
 PROOF: $\langle 2 \rangle 3$
 $\langle 2 \rangle 7$. $y \neq z$
 PROOF: $\langle 2 \rangle 3$
 $\langle 2 \rangle 8$. $x \leq z$
 PROOF: From $\langle 2 \rangle 4$ and $\langle 2 \rangle 6$ since \leq is transitive.
 $\langle 2 \rangle 9$. $x \neq z$
 $\langle 3 \rangle 1$. ASSUME: for a contradiction $x = z$
 $\langle 3 \rangle 2$. $y \leq x$
 PROOF: $\langle 2 \rangle 6$, $\langle 3 \rangle 1$
 $\langle 3 \rangle 3$. $x = y$
 PROOF: From $\langle 2 \rangle 4$ and $\langle 3 \rangle 2$ since \leq is antisymmetric.
 $\langle 3 \rangle 4$. Q.E.D.
 PROOF: $\langle 2 \rangle 5$ and $\langle 3 \rangle 3$ form a contradiction.
 $\langle 2 \rangle 10$. $x < z$
 PROOF: $\langle 2 \rangle 8$, $\langle 2 \rangle 9$
 $\langle 1 \rangle 5$. If there exists a linear order \leq on X such that $\forall x, y \in X (x < y \Leftrightarrow x \leq y \wedge x \neq y)$, then $\forall x, y \in X (x < y \vee x = y \vee y < x)$.
 $\langle 2 \rangle 1$. LET: \leq be a linear order on X .
 $\langle 2 \rangle 2$. ASSUME: $\forall x, y \in X (x < y \Leftrightarrow x \leq y \wedge x \neq y)$
 $\langle 2 \rangle 3$. LET: $x, y \in X$
 $\langle 2 \rangle 4$. $x \leq y$ or $y \leq x$
 $\langle 2 \rangle 5$. CASE: $x \leq y$
 PROOF: Then $x < y$ or $x = y$ using $\langle 2 \rangle 2$.
 $\langle 2 \rangle 6$. CASE: $y \leq x$
 PROOF: Then $y < x$ or $x = y$ using $\langle 2 \rangle 2$.
 $\langle 1 \rangle 6$. If there exists a linear order \leq on X such that $\forall x, y \in X (x < y \Leftrightarrow x \leq y \wedge x \neq y)$, then $\forall x, y \in X (x \leq y \Leftrightarrow x < y \vee x = y)$.
 PROOF: Trivial.
 $\langle 1 \rangle 7$. If $<$ is irreflexive, transitive and satisfies $\forall x, y \in X (x < y \vee x = y \vee y < x)$, then the relation \leq defined by $\forall x, y \in X (x \leq y \Leftrightarrow x < y \vee x = y)$ is a linear order on X .
 $\langle 2 \rangle 1$. ASSUME: $<$ is irreflexive.

$\langle 2 \rangle 2$. ASSUME: $<$ is transitive.
 $\langle 2 \rangle 3$. ASSUME: $\forall x, y \in X (x < y \vee x = y \vee y < x)$
 $\langle 2 \rangle 4$. LET: \leq be the relation defined by $\forall x, y (x \leq y \Leftrightarrow x < y \vee x = y)$
 $\langle 2 \rangle 5$. \leq is reflexive.
 PROOF: Immediate from $\langle 2 \rangle 4$.
 $\langle 2 \rangle 6$. \leq is transitive.
 $\langle 3 \rangle 1$. LET: $x, y, z \in X$
 $\langle 3 \rangle 2$. ASSUME: $x \leq y$ and $y \leq z$
 $\langle 3 \rangle 3$. $x < y$ or $x = y$
 PROOF: $\langle 2 \rangle 4, \langle 3 \rangle 2$
 $\langle 3 \rangle 4$. $y < z$ or $y = z$
 PROOF: $\langle 2 \rangle 4, \langle 3 \rangle 2$
 $\langle 3 \rangle 5$. CASE: $x < y$ and $y < z$
 $\langle 4 \rangle 1$. $x < z$
 PROOF: $\langle 2 \rangle 2$
 $\langle 4 \rangle 2$. $x \leq z$
 PROOF: $\langle 2 \rangle 4, \langle 4 \rangle 1$
 $\langle 3 \rangle 6$. CASE: $x = y$
 PROOF: Then $x \leq z$ from $\langle 3 \rangle 2$
 $\langle 3 \rangle 7$. CASE: $y = z$
 PROOF: Then $x \leq z$ from $\langle 3 \rangle 2$
 $\langle 2 \rangle 7$. \leq is antisymmetric.
 $\langle 3 \rangle 1$. LET: $x, y \in X$
 $\langle 3 \rangle 2$. ASSUME: $x \leq y$
 $\langle 3 \rangle 3$. ASSUME: $y \leq x$
 $\langle 3 \rangle 4$. $x < y$ or $x = y$
 PROOF: $\langle 2 \rangle 4, \langle 3 \rangle 2$
 $\langle 3 \rangle 5$. $y < x$ or $y = x$
 PROOF: $\langle 2 \rangle 4, \langle 3 \rangle 2$
 $\langle 3 \rangle 6$. $\neg(x < y \wedge y < x)$
 $\langle 4 \rangle 1$. ASSUME: for a contradiction $x < y$ and $y < x$
 $\langle 4 \rangle 2$. $x < x$
 PROOF: $\langle 2 \rangle 2, \langle 4 \rangle 1$
 $\langle 4 \rangle 3$. Q.E.D.
 PROOF: $\langle 2 \rangle 1$ and $\langle 4 \rangle 2$ form a contradiction.
 $\langle 3 \rangle 7$. $x = y$
 PROOF: $\langle 3 \rangle 4, \langle 3 \rangle 5, \langle 3 \rangle 6$
 $\langle 2 \rangle 8$. $\forall x, y \in X (x \leq y \vee y \leq x)$
 $\langle 3 \rangle 1$. LET: $x, y \in X$
 $\langle 3 \rangle 2$. $x < y$ or $x = y$ or $y < x$
 PROOF: $\langle 2 \rangle 3, \langle 3 \rangle 1$
 $\langle 3 \rangle 3$. $x \leq y$ or $y \leq x$
 PROOF: $\langle 2 \rangle 4, \langle 3 \rangle 2$

□

Theorem 5.21 (Zorn's Lemma). *Let X be a poset such that every chain in X*

has an upper bound. Then X has a maximal element.

PROOF:

⟨1⟩1. PICK a choice function f for X .

⟨1⟩2. LET: \mathcal{X} be the set of chains in X .

⟨1⟩3. For all $A \in \mathcal{X}$,

LET: $\hat{A} = \{x \in X : A \cup \{x\} \in \mathcal{X}\}$

⟨1⟩4. LET: $g : \mathcal{X} \rightarrow \mathcal{X}$ be the function

$$g(A) = \begin{cases} A \cup \{f(\hat{A} - A)\} & \text{if } \hat{A} - A \neq \emptyset \\ A & \text{if } \hat{A} - A = \emptyset \end{cases}$$

⟨1⟩5. For $\mathcal{T} \subseteq \mathcal{X}$, let us say \mathcal{T} is a *tower* iff:

- $\emptyset \in \mathcal{T}$
- $\forall A \in \mathcal{T}. g(A) \in \mathcal{T}$
- For every chain \mathcal{C} in \mathcal{T} , we have $\bigcup \mathcal{C} \in \mathcal{T}$

⟨1⟩6. LET: \mathcal{T}_0 be the intersection of the set of all towers.

PROOF: The set of all towers is nonempty since \mathcal{X} is a tower.

⟨1⟩7. LET: $A = \bigcup \mathcal{T}_0$

⟨1⟩8. A is a chain in X .

⟨2⟩1. \mathcal{T}_0 is a chain under \subseteq

⟨3⟩1. Given $C \in \mathcal{T}_0$, let us say that C is *comparable* iff, for all $A \in \mathcal{T}_0$, either $A \subseteq C$ or $C \subseteq A$.

⟨3⟩2. For all $A, C \in \mathcal{T}_0$, if C is comparable and $A \subsetneq C$ then $g(A) \subseteq C$.

PROOF: Since $g(A) - A$ has at most one element, so if $A \subsetneq C \subseteq g(A)$ then $C = g(A)$.

⟨3⟩3. For $C \in \mathcal{T}_0$ comparable,

LET: $\mathcal{U}_C = \{A \in \mathcal{T}_0 : A \subseteq C \vee g(C) \subseteq A\}$

⟨3⟩4. For $C \in \mathcal{T}_0$ comparable, \mathcal{U}_C is a tower.

⟨4⟩1. LET: $C \in \mathcal{T}_0$ be comparable

⟨4⟩2. $\emptyset \in \mathcal{U}_C$

PROOF: Since $\emptyset \subseteq C$.

⟨4⟩3. $\forall A \in \mathcal{U}_C. g(A) \in \mathcal{U}_C$

PROOF: By ⟨1⟩8.

⟨4⟩4. For every chain $\mathcal{C} \subseteq \mathcal{U}_C$ we have $\bigcup \mathcal{C} \in \mathcal{U}_C$

⟨5⟩1. LET: $\mathcal{C} \subseteq \mathcal{U}_C$ be a chain.

⟨5⟩2. CASE: $\exists A \in \mathcal{C}. g(C) \subseteq A$

PROOF: Then $g(C) \subseteq \bigcup \mathcal{C}$

⟨5⟩3. CASE: $\forall A \in \mathcal{C}. A \subseteq C$

PROOF: Then $\bigcup \mathcal{C} \subseteq C$.

⟨3⟩5. For $C \in \mathcal{T}_0$ comparable, $\mathcal{U}_C = \mathcal{T}_0$.

⟨3⟩6. For $C \in \mathcal{T}_0$ comparable we have $g(C)$ is comparable.

PROOF: Since for all $A \in \mathcal{T}_0$ either $A \subseteq C \subseteq g(C)$ or $g(C) \subseteq A$.

⟨3⟩7. The set of comparable sets in \mathcal{T}_0 is a tower.

⟨4⟩1. \emptyset is comparable.

PROOF: $\forall A \in \mathcal{T}_0. \emptyset \subseteq A$

$\langle 4 \rangle 2$. For all $C \in \mathcal{T}_0$, if A is comparable then $g(C)$ is comparable.
 PROOF: $\langle 3 \rangle 6$
 $\langle 4 \rangle 3$. For every chain $\mathcal{C} \subseteq \mathcal{T}_0$ of comparable sets, we have $\bigcup \mathcal{C}$ is comparable.
 $\langle 5 \rangle 1$. LET: $\mathcal{C} \subseteq \mathcal{T}_0$ be a chain of comparable sets.
 $\langle 5 \rangle 2$. LET: $A \in \mathcal{T}_0$
 $\langle 5 \rangle 3$. CASE: there exists $C \in \mathcal{C}$ such that $A \subseteq C$
 PROOF: Then $A \subseteq \bigcup \mathcal{C}$.
 $\langle 5 \rangle 4$. CASE: for all $C \in \mathcal{C}$ we have $C \subseteq A$
 PROOF: Then $\bigcup \mathcal{C} \subseteq A$.
 $\langle 3 \rangle 8$. Every set in \mathcal{T}_0 is comparable.
 $\langle 2 \rangle 2$. LET: $x, y \in A$
 $\langle 2 \rangle 3$. PICK $A, C \in \mathcal{T}_0$ such that $x \in A$ and $y \in C$
 $\langle 2 \rangle 4$. ASSUME: w.l.o.g. $A \subseteq C$
 $\langle 2 \rangle 5$. $x, y \in C$
 $\langle 2 \rangle 6$. $x \leq y$ or $y \leq x$
 PROOF: Since $C \in \mathcal{X}$ so C is a chain.
 $\langle 1 \rangle 9$. PICK an upper bound u for A .
 $\langle 1 \rangle 10$. $A \in \mathcal{T}_0$
 PROOF: Since \mathcal{T}_0 is a chain in \mathcal{T}_0 so $\bigcup \mathcal{T}_0 \in \mathcal{T}_0$.
 $\langle 1 \rangle 11$. $g(A) \in \mathcal{T}_0$
 $\langle 1 \rangle 12$. $g(A) \subseteq A$
 $\langle 1 \rangle 13$. $g(A) = A$
 $\langle 1 \rangle 14$. $\hat{A} - A = \emptyset$
 $\langle 1 \rangle 15$. $u \in A$
 PROOF: Since $A \cup \{u\}$ is a chain so $u \in \hat{A}$ and therefore $u \in A$.
 $\langle 1 \rangle 16$. u is maximal in X .
 $\langle 2 \rangle 1$. LET: $x \in X$
 $\langle 2 \rangle 2$. ASSUME: $u \leq x$
 $\langle 2 \rangle 3$. $A \cup \{x\}$ is a chain.
 $\langle 2 \rangle 4$. $x \in A$
 $\langle 2 \rangle 5$. $x \leq u$
 $\langle 2 \rangle 6$. $x = u$

□

Definition 5.22 (Cofinal). Let X be a poset and $A \subseteq X$. Then A is *cofinal* iff, for all $x \in X$, there exists $a \in A$ such that $x \leq a$.

Definition 5.23 (Order Isomorphism). Two posets X and Y are (*order*) *isomorphic*, $X \cong Y$ iff there exists an order preserving one-to-one correspondence f between them. We write $f : X \cong Y$ and call f a (*order*) *isomorphism*.

Proposition 5.24. Let X and Y be posets. Let f be a one-to-one correspondence between X and Y . Then f is a similarity if and only if, for all $x, y \in X$, we have $x < y$ iff $f(x) < f(y)$.

PROOF: Easy. □

Proposition 5.25. *For any poset X we have $I_X : X \cong X$.*

PROOF: Easy. \square

Proposition 5.26. *If $f : X \cong Y$ then $f^{-1} : Y \cong X$.*

PROOF: Easy. \square

Proposition 5.27. *If $f : X \cong Y$ and $g : Y \cong Z$ then $g \circ f : X \cong Z$.*

PROOF: Easy. \square

Corollary 5.27.1. *For any set E , similarity is an equivalence relation on the set of all posets that are subsets of E .*

Definition 5.28 (Open Interval). Let X be a linearly ordered set. Let $a, b \in X$ with $a < b$. The *open interval* (a, b) is

$$(a, b) := \{x \in X : a < x < b\} .$$

Definition 5.29 (Lexicographical Order). Let A and B be linearly ordered sets. The *lexicographical order* on $A \times B$ is the relation \leq defined by $(a, b) \leq (x, y)$ iff $a < x$ or $(a = x \text{ and } b \leq y)$.

Definition 5.30 (Least Upper Bound Property). A linearly ordered set A has the *least upper bound property* iff every nonempty subset of A bounded above has a least upper bound.

Proposition 5.31. *Let A be a linearly ordered set. Then A has the least upper bound property if and only if every nonempty subset of A bounded below has a greatest lower bound.*

PROOF:

$\langle 1 \rangle 1$. For any linearly ordered set A , if A has the least upper bound property then every nonempty subset of A bounded below has a greatest lower bound.

$\langle 2 \rangle 1$. LET: A be a linearly ordered set.

$\langle 2 \rangle 2$. ASSUME: A has the least upper bound property.

$\langle 2 \rangle 3$. LET: S be a nonempty subset of A bounded below.

$\langle 2 \rangle 4$. LET: $S \downarrow$ be the set of all lower bounds for S .

$\langle 2 \rangle 5$. $S \downarrow \neq \emptyset$

PROOF: Since S is bounded below by $\langle 2 \rangle 3$.

$\langle 2 \rangle 6$. $S \downarrow$ is bounded above.

$\langle 3 \rangle 1$. PICK $s \in S$

PROOF: S is nonempty by $\langle 2 \rangle 3$.

$\langle 3 \rangle 2$. s is an upper bound for $S \downarrow$.

$\langle 2 \rangle 7$. LET: l be the supremum of $S \downarrow$.

PROOF: $\langle 2 \rangle 2, \langle 2 \rangle 5, \langle 2 \rangle 6$

$\langle 2 \rangle 8$. l is a lower bound for S .

$\langle 3 \rangle 1$. LET: $s \in S$

$\langle 3 \rangle 2$. s is an upper bound for $S \downarrow$

$\langle 3 \rangle 3$. $l \leq s$

PROOF: $\langle 2 \rangle 7$, $\langle 3 \rangle 2$

$\langle 2 \rangle 9$. For any lower bound l' for S we have $l' \leq l$.

PROOF: Since l is an upper bound for $S \downarrow$ by $\langle 2 \rangle 7$.

$\langle 1 \rangle 2$. For any linearly ordered set A , if every nonempty subset of A bounded below has a greatest lower bound then A has the least upper bound property.

$\langle 2 \rangle 1$. LET: (A, \leq) be a linearly ordered set.

$\langle 2 \rangle 2$. ASSUME: Every nonempty subset of A bounded below has a greatest lower bound.

$\langle 2 \rangle 3$. (A, \geq) has the least upper bound property.

$\langle 2 \rangle 4$. Every nonempty subset of A bounded below with respect to \geq has a greatest lower bound with respect to \geq .

$\langle 2 \rangle 5$. Every nonempty subset of A bounded above with respect to \leq has a least upper bound with respect to \leq .

□

5.3 Linear Continua

Definition 5.32 (Linear Continuum). A *linear continuum* is a linearly ordered set X with the least upper bound property such that, for all $x, y \in X$, if $x < y$ then there exists $z \in X$ such that $x < z < y$.

5.4 Well Orderings

Definition 5.33 (Well Ordered Set). A poset X is *well ordered*, and its ordering is a *well ordering*, iff every nonempty subset of X has a least element.

Proposition 5.34. Every well ordered set is totally ordered.

PROOF: For all x and y we have $\{x, y\}$ has a least element, so $x \leq y$ or $y \leq x$. □

Theorem 5.35 (Transfinite Induction). Let X be a well ordered set. Let $S \subseteq X$ satisfy:

$$\forall x \in X (\forall y < x. y \in S) \Rightarrow x \in S .$$

Then $S = X$.

PROOF: We have $X - S$ has no least element, so $X - S = \emptyset$. □

Definition 5.36 (Continuation). Let A and B be well ordered sets. Then B is a *continuation* of A iff there exists $b \in B$ such that $A = s(b)$ and the order on A is the restriction of the order on B to A .

Proposition 5.37. Let \mathcal{C} be a set of well ordered sets that is totally ordered under continuation. Then there exists a unique well ordering on $\bigcup \mathcal{C}$ such that $\bigcup \mathcal{C}$ is a continuation of every element of \mathcal{C} .

PROOF: Define \leq on $\bigcup \mathcal{C}$ by: $x \leq y$ iff there exists $C \in \mathcal{C}$ such that $x, y \in C$ and $x \leq y$ in C . \square

Proposition 5.38. *Every totally ordered set has a cofinal well ordered subset.*

PROOF:

$\langle 1 \rangle 1$. LET: X be a totally ordered set.

$\langle 1 \rangle 2$. LET: \mathcal{C} be the poset of all well ordered subsets of X under continuation.

$\langle 1 \rangle 3$. Every chain in \mathcal{C} has an upper bound.

PROOF: Proposition 5.37.

$\langle 1 \rangle 4$. PICK a maximal element C of \mathcal{C}

PROVE: C is cofinal

PROOF: Zorn's Lemma

$\langle 1 \rangle 5$. LET: $x \in X$

$\langle 1 \rangle 6$. We cannot have $\forall c \in C. c < x$

PROOF: Then $C \cup \{x\}$ would be a larger chain.

$\langle 1 \rangle 7$. $\exists c \in C. x \leq c$

\square

Theorem 5.39 (Well Ordering Theorem). *Every set can be well ordered.*

PROOF:

$\langle 1 \rangle 1$. LET: X be a set.

$\langle 1 \rangle 2$. LET: \mathcal{W} be the poset of all well ordered subsets of X under continuation.

$\langle 1 \rangle 3$. Every chain in \mathcal{W} has an upper bound.

PROOF: Proposition 5.37.

$\langle 1 \rangle 4$. PICK a maximal $M \in \mathcal{W}$

PROOF: Zorn's Lemma

$\langle 1 \rangle 5$. $M = X$

PROOF: If $x \in X - M$ then $M \cup \{x\}$ with x as the greatest element is a continuation of M .

\square

Theorem 5.40 (Transfinite Recursion). *Let W be a well ordered set and X a set. Let S be the set of all functions f such that $\text{ran } f \subseteq X$, and there exists $a \in W$ such that $\text{dom } f = s(a)$. Then there exists a unique function $U : W \rightarrow X$ such that*

$$\forall a \in W. U(a) = f(U \upharpoonright s(a)) .$$

PROOF:

$\langle 1 \rangle 1$. Let us say that a subset $A \subseteq W \times X$ is *f-closed* iff, whenever $a \in W$ and $t : s(a) \rightarrow X$ satisfies $\forall c < a. (c, t(c)) \in A$, then $(a, f(t)) \in A$.

$\langle 1 \rangle 2$. LET: U be the intersection of the set of *f-closed* subsets of $W \times X$

PROOF: This set is nonempty since $W \times X$ is *f-closed*.

$\langle 1 \rangle 3$. U is *f-closed*.

$\langle 1 \rangle 4$. U is a function.

$\langle 2 \rangle 1$. LET: $P(a)$ be the property: there is at most one $x \in X$ such that $(a, x) \in U$

$\langle 2 \rangle 2$. LET: $a \in W$
 $\langle 2 \rangle 3$. ASSUME: as transfinite induction hypothesis $\forall c < a. P(c)$
 $\langle 2 \rangle 4$. LET: $(a, x), (a, y) \in U$
 $\langle 2 \rangle 5$. $x = f(U \upharpoonright c)$
 PROOF: If not then $U - \{(a, x)\}$ would be f -closed.
 $\langle 2 \rangle 6$. $y = f(U \upharpoonright c)$
 $\langle 2 \rangle 7$. $x = y$
 $\langle 1 \rangle 5$. $\text{dom } U = W$
 $\langle 2 \rangle 1$. LET: $a \in W$
 $\langle 2 \rangle 2$. ASSUME: as transfinite induction hypothesis $\forall c < a. c \in \text{dom } U$
 $\langle 2 \rangle 3$. $(a, f(U \upharpoonright s(a))) \in U$
 $\langle 1 \rangle 6$. If $U' : W \rightarrow X$ and $\forall a \in W. U'(a) = f(U' \upharpoonright s(a))$, then $U' = U$.
 PROOF: Prove $U'(a) = U(a)$ by transfinite induction on a .
 \square

Proposition 5.41. *Let X be a well ordered set and f a similarity between X and a subset of X . Then, for all $a \in X$, we have $a \leq f(a)$.*

PROOF:
 $\langle 1 \rangle 1$. LET: $a \in X$
 $\langle 1 \rangle 2$. ASSUME: as transfinite induction hypothesis $\forall c < a. c \leq f(c)$
 $\langle 1 \rangle 3$. ASSUME: for a contradiction $f(a) < a$
 $\langle 1 \rangle 4$. $f(a) \leq f(f(a))$
 PROOF: $\langle 1 \rangle 2$
 $\langle 1 \rangle 5$. $f(f(a)) < f(a)$
 PROOF: From $\langle 1 \rangle 3$ since f is a similarity.
 $\langle 1 \rangle 6$. Q.E.D.
 PROOF: This is a contradiction.
 \square

Proposition 5.42. *Let X and Y be well ordered sets. Then there is at most one similarity between them.*

PROOF:
 $\langle 1 \rangle 1$. LET: $f, g : X \cong Y$
 PROVE: $\forall a \in X. f(a) = g(a)$
 $\langle 1 \rangle 2$. LET: $a \in X$
 $\langle 1 \rangle 3$. ASSUME: as transfinite induction hypothesis $\forall c < a. f(c) = g(c)$
 $\langle 1 \rangle 4$. $f(a)$ is the least element of $Y - \{f(c) : c < a\}$
 $\langle 1 \rangle 5$. $g(a)$ is the least element of $Y - \{g(c) : c < a\}$
 $\langle 1 \rangle 6$. $f(a) = g(a)$
 \square

Proposition 5.43. *A well ordered set is not similar to any of its initial segments.*

PROOF:
 $\langle 1 \rangle 1$. LET: X be a well ordered set.

⟨1⟩2. ASSUME: for a contradiction $f : X \cong s(a)$ for some $a \in X$

⟨1⟩3. $f(a) < a$

⟨1⟩4. Q.E.D.

PROOF: This contradicts Proposition 5.41.

□

Theorem 5.44 (Comparability Theorem). *Given well ordered sets X and Y , either $X \cong Y$, or X is similar to an initial segment of Y , or Y is similar to an initial segment of X .*

PROOF:

⟨1⟩1. LET: $X_0 = \{a \in X : \exists b \in Y. s(a) \cong s(b)\}$

⟨1⟩2. LET: $U : X_0 \rightarrow Y$ be the function: for $a \in X_0$, we have $U(a)$ is the unique element in Y such that $s(a) \cong s(U(a))$

⟨1⟩3. LET: $Y_0 = \text{ran } U$

⟨1⟩4. Either $X_0 = X$ or there exists $a \in X$ such that $X_0 = s(a)$

⟨2⟩1. ASSUME: $X_0 \neq X$

⟨2⟩2. LET: a be the least element of $X - X_0$

⟨2⟩3. LET: $x \in X_0$

PROVE: $x < a$

⟨2⟩4. PICK $f : s(x) \cong s(U(x))$

⟨2⟩5. ASSUME: for a contradiction $a < x$

⟨2⟩6. $f \upharpoonright s(a) : s(a) \cong s(f(a))$

⟨2⟩7. $a \in X_0$

⟨2⟩8. Q.E.D.

PROOF: This is a contradiction.

⟨1⟩5. Either $Y_0 = Y$ or there exists $b \in Y$ such that $Y_0 = s(b)$

PROOF: Similar.

⟨1⟩6. CASE: $X_0 = X$ and $Y_0 = Y$

PROOF: Then $U : X \cong Y$.

⟨1⟩7. CASE: $X_0 = X$ and $Y_0 \neq Y$

PROOF: Then $U : X \cong s(b)$ where $Y_0 = s(b)$.

⟨1⟩8. CASE: $X_0 \neq X$ and $Y_0 = Y$

PROOF: Then $U : s(a) \cong Y$ where $X_0 = s(a)$.

⟨1⟩9. CASE: $X_0 \neq X$ and $Y_0 \neq Y$

⟨2⟩1. LET: $X_0 = s(a)$ and $Y_0 = s(b)$

⟨2⟩2. $U : s(a) \cong s(b)$

⟨2⟩3. $a \in X_0$

⟨2⟩4. Q.E.D.

PROOF: This is a contradiction.

□

Corollary 5.44.1. *Let X be a well ordered set. Then any subset A of X is either similar to X or to an initial segment of X .*

PROOF: We cannot have X is similar to an initial segment of A , say $f : X \cong \{x \in A : x < a\}$, because then we would have $f(a) < a$ contradicting Proposition 5.41. □

Corollary 5.44.2. *For any sets X and Y , either there exists an injective function $X \rightarrow Y$, or there exists an injective function $Y \rightarrow X$.*

PROOF: Using the Well Ordering Theorem. \square

Chapter 6

Natural Numbers

6.1 Natural Numbers

Definition 6.1 (Successor). The *successor* of a set x , x^+ , is defined by

$$x^+ := x \cup \{x\} .$$

Definition 6.2. We define

$$0 = \emptyset$$

$$1 = 0^+$$

$$2 = 1^+$$

etc.

Definition 6.3 (Characteristic Function). Let X be a set and $A \subseteq X$. The *characteristic function* of A is the function $\chi_A : X \rightarrow 2$ defined by

$$\chi_A(x) = \begin{cases} 1 & \text{if } x \in A \\ 0 & \text{if } x \notin A \end{cases}$$

Theorem 6.4. Let X be a set. The function $\chi : \mathcal{P}X \rightarrow 2^X$ that maps a subset A of X to χ_A is a one-to-one correspondence.

PROOF: Easy. \square

Definition 6.5. The set ω of *natural numbers* is the set such that:

- $0 \in \omega$
- For all $n \in \omega$ we have $n^+ \in \omega$
- For any set X , if $0 \in X$ and $\forall n \in X. n^+ \in X$ then $\omega \subseteq X$

PROOF: To show this exists, pick a set A such that $0 \in A$ and $\forall n \in A. n^+ \in A$ (by the Axiom of Infinity), and let $\omega = \bigcap \{X \in \mathcal{P}A : 0 \in X \wedge \forall n \in X. n^+ \in X\}$.
 \square

Definition 6.6 (Sequence). A *finite sequence* is a family whose index set is a natural number. An *infinite sequence* is a family whose index set is ω .

Given a finite sequence of sets $\{A_i\}_{i \in n^+}$, we write $\bigcup_{i=0}^n A_i$ for $\bigcup_{i \in n^+} A_i$. Given an infinite sequence of sets $\{A_i\}_{i \in \omega}$, we write $\bigcup_{i=0}^{\infty} A_i$ for $\bigcup_{i \in \omega} A_i$.

We make similar definitions for \bigcap and \times .

Proposition 6.7. For any natural numbers m and n , if $m \in n$ then $m^+ \in n^+$.

PROOF:

$\langle 1 \rangle 1$. LET: $P(n)$ be the property $\forall m \in n. m^+ \in n^+$

$\langle 1 \rangle 2$. $P(0)$

PROOF: Vacuous.

$\langle 1 \rangle 3$. For any natural number n , if $P(n)$ then $P(n^+)$.

$\langle 2 \rangle 1$. LET: n be a natural number.

$\langle 2 \rangle 2$. ASSUME: $P(n)$

$\langle 2 \rangle 3$. LET: $m \in n^+$

$\langle 2 \rangle 4$. $m \in n$ or $m = n$

$\langle 2 \rangle 5$. $m^+ \in n^+$ or $m^+ = n^+$

PROOF: $\langle 2 \rangle 2$

$\langle 2 \rangle 6$. CASE: $m^+ \in n^{++}$

\square

Theorem 6.8 (Principle of Mathematical Induction). For any subset S of ω , if $0 \in S$ and $\forall n \in S. n^+ \in S$, then $S = \omega$.

PROOF: From the definition of ω . \square

Proposition 6.9.

$$\forall n \in \omega. \forall x \in n. n \not\subseteq x$$

PROOF:

$\langle 1 \rangle 1$. $\forall x \in 0. 0 \not\subseteq x$

PROOF: Vacuous.

$\langle 1 \rangle 2$. For any natural number n , if $\forall x \in n. n \not\subseteq x$ then $\forall x \in n^+. n^+ \not\subseteq x$.

$\langle 2 \rangle 1$. LET: n be a natural number.

$\langle 2 \rangle 2$. ASSUME: $\forall x \in n. n \not\subseteq x$

$\langle 2 \rangle 3$. LET: $x \in n^+$

$\langle 2 \rangle 4$. ASSUME: for a contradiction $n^+ \subseteq x$

$\langle 2 \rangle 5$. $x \in n$ or $x = n$

$\langle 2 \rangle 6$. CASE: $x \in n$

PROOF: Then we have $n \subseteq n^+ \subseteq x$ contradicting $\langle 2 \rangle 2$.

$\langle 2 \rangle 7$. CASE: $x = n$

PROOF: Then we have $n \in n^+ \subseteq x = n$ and $n \subseteq n$ contradicting $\langle 2 \rangle 2$.

\square

Corollary 6.9.1. *For any natural number n we have $n \notin n$.*

Corollary 6.9.2. *For any natural number n we have $n \neq n^+$.*

Definition 6.10 (Transitive Set). A set E is a *transitive set* iff, whenever $x \in y \in E$, then $x \in E$.

Proposition 6.11. *Every natural number is a transitive set.*

PROOF:

$\langle 1 \rangle 1$. 0 is a transitive set.

PROOF: Vacuously, if $x \in y \in 0$ then $x \in 0$.

$\langle 1 \rangle 2$. For any natural number n , if n is a transitive set, then n^+ is a transitive set.

$\langle 2 \rangle 1$. LET: n be a natural number.

$\langle 2 \rangle 2$. ASSUME: n is a transitive set.

$\langle 2 \rangle 3$. LET: $x \in y \in n^+$

$\langle 2 \rangle 4$. $y \in n$ or $y = n$

$\langle 2 \rangle 5$. CASE: $y \in n$

$\langle 3 \rangle 1$. $x \in n$

PROOF: $\langle 2 \rangle 2$, $\langle 2 \rangle 3$, $\langle 2 \rangle 5$.

$\langle 3 \rangle 2$. $x \in n^+$

$\langle 2 \rangle 6$. CASE: $y = n$

$\langle 3 \rangle 1$. $x \in n$

PROOF: $\langle 2 \rangle 3$, $\langle 2 \rangle 6$

$\langle 3 \rangle 2$. $x \in n^+$

□

Proposition 6.12. *For any natural numbers m and n , if $m^+ = n^+$ then $m = n$.*

PROOF:

$\langle 1 \rangle 1$. LET: m and n be natural numbers.

$\langle 1 \rangle 2$. ASSUME: $m^+ = n^+$

$\langle 1 \rangle 3$. $m \in m^+ = n^+$

$\langle 1 \rangle 4$. $m \in n$ or $m = n$

$\langle 1 \rangle 5$. $n \in n^+ = m^+$

$\langle 1 \rangle 6$. $n \in m$ or $n = m$

$\langle 1 \rangle 7$. We cannot have $m \in n$ and $n \in m$

$\langle 2 \rangle 1$. ASSUME: for a contradiction $m \in n$ and $n \in m$

$\langle 2 \rangle 2$. $m \in m$

PROOF: Since m is a transitive set (Proposition 6.11).

$\langle 2 \rangle 3$. Q.E.D.

PROOF: This contradicts Proposition 6.9.

$\langle 1 \rangle 8$. $m = n$

□

Theorem 6.13 (Recursion Theorem). *Let X be a set. Let $a \in X$. Let $f : X \rightarrow X$. There exists a function $u : \omega \rightarrow X$ such that $u(0) = a$ and, for all $n \in \omega$, we have $u(n^+) = f(u(n))$.*

PROOF:

$\langle 1 \rangle 1$. LET: $\mathcal{C} = \{A \in \mathcal{P}(\omega \times X) : (0, a) \in A \wedge \forall n \in \omega. \forall x \in X. (n, x) \in A \Rightarrow (n^+, f(x)) \in A\}$

$\langle 1 \rangle 2$. $\mathcal{C} \neq \emptyset$

PROOF: $\omega \times X \in \mathcal{C}$

$\langle 1 \rangle 3$. LET: $u = \bigcap \mathcal{C}$

$\langle 1 \rangle 4$. $u \in \mathcal{C}$

$\langle 1 \rangle 5$. u is a function.

$\langle 2 \rangle 1$. LET: $P(n)$ be the property: $\forall x, y \in X. (n, x) \in u \wedge (n, y) \in u \Rightarrow x = y$

$\langle 2 \rangle 2$. $P(0)$

$\langle 3 \rangle 1$. $\forall x \in X. (0, x) \in u \Rightarrow x = a$

PROOF: If $(0, x) \in u$ and $x \neq a$ then $u - \{(0, x)\} \in \mathcal{C}$ and so $u - \{(0, x)\} \subseteq u$, which is impossible.

$\langle 2 \rangle 3$. For every natural number n , if $P(n)$ then $P(n^+)$.

$\langle 3 \rangle 1$. LET: n be a natural number.

$\langle 3 \rangle 2$. ASSUME: $P(n)$

$\langle 3 \rangle 3$. LET: $x, y \in X$

$\langle 3 \rangle 4$. ASSUME: $(n^+, x), (n^+, y) \in u$

$\langle 3 \rangle 5$. PICK $x', y' \in X$ such that $(n, x') \in u$, $(n, y') \in u$ and $f(x') = x$ and $f(y') = y$

PROOF: If no such x' exists then $u - \{(n^+, x)\} \in \mathcal{C}$ and so $u - \{(n^+, x)\} \subseteq u$ which is impossible. Similarly for y' .

$\langle 3 \rangle 6$. $x' = y'$

PROOF: $\langle 3 \rangle 2$

$\langle 3 \rangle 7$. $x = y$

□

Proposition 6.14. *For any natural number n , either $n = 0$ or there exists a natural number m such that $n = m^+$.*

PROOF: Easy induction on n . □

Proposition 6.15. *ω is a transitive set.*

PROOF:

$\langle 1 \rangle 1$. LET: $P(n)$ be the property $\forall x \in n. x \in \omega$

$\langle 1 \rangle 2$. $P(0)$

PROOF: Vacuous.

$\langle 1 \rangle 3$. For any natural number n , if $P(n)$ then $P(n^+)$.

$\langle 2 \rangle 1$. LET: n be a natural number.

$\langle 2 \rangle 2$. ASSUME: $P(n)$

$\langle 2 \rangle 3$. LET: $x \in n^+$

$\langle 2 \rangle 4$. $x \in n$ or $x = n$

$\langle 2 \rangle 5$. CASE: $x \in n$

PROOF: Then $x \in \omega$ by $\langle 2 \rangle 2$.

$\langle 2 \rangle 6$. CASE: $x = n$

PROOF: Then $x \in \omega$ by $\langle 2 \rangle 1$.

□

Proposition 6.16. *For any natural number n and any nonempty subset $E \subseteq n$, there exists $k \in E$ such that $\forall m \in E. k = m \vee k \in m$.*

PROOF:

⟨1⟩1. LET: $P(n)$ be the property: for any nonempty subset $E \subseteq n$, there exists $k \in E$ such that $\forall m \in E. k = m \vee k \in m$

⟨1⟩2. $P(0)$

PROOF: Vacuous as there is no nonempty subset of 0.

⟨1⟩3. For any natural number n , if $P(n)$ then $P(n^+)$.

⟨2⟩1. LET: n be a natural number.

⟨2⟩2. ASSUME: $P(n)$

⟨2⟩3. LET: E be a nonempty subset of n^+

⟨2⟩4. CASE: $E - \{n\} = \emptyset$

PROOF: Then $E = \{n\}$ so take $k = n$.

⟨2⟩5. CASE: $E - \{n\} \neq \emptyset$

⟨3⟩1. PICK $k \in E - \{n\}$ such that $\forall m \in E - \{n\}. k = m \vee k \in m$

PROOF: By ⟨2⟩2.

⟨3⟩2. $\forall m \in E. k = m \vee k \in m$

PROOF: Since $k \in n$.

□

Chapter 7

Ordinal Numbers

Definition 7.1 (Ordinal (Number)). An *ordinal (number)* is a well ordered set α such that $\forall \xi \in \alpha. s(\xi) = \xi$.

Given ordinals α, β , we write $\alpha < \beta$ iff $\alpha \in \beta$.

Proposition 7.2. *Every natural number is an ordinal.*

PROOF: Easy. \square

Proposition 7.3. ω is an ordinal.

PROOF: Easy. \square

Proposition 7.4. If α is an ordinal number then so is α^+ .

PROOF: Easy. \square

Proposition 7.5. Let α be an ordinal and $\eta, \xi \in \alpha$. Then $\eta < \xi$ if and only if $\eta \in \xi$.

PROOF: Easy. \square

Proposition 7.6. Every ordinal is a transitive set.

PROOF: Easy. \square

Proposition 7.7. Every element of an ordinal is an ordinal.

PROOF: Easy. \square

Proposition 7.8. Similar ordinals are equal.

PROOF:

$\langle 1 \rangle 1$. LET: α, β be ordinals.

$\langle 1 \rangle 2$. LET: $f : \alpha \cong \beta$ be a similarity.

PROVE: $\forall \xi \in \alpha. f(\xi) = \xi$

$\langle 1 \rangle 3$. LET: $\xi \in \alpha$

$\langle 1 \rangle 4$. ASSUME: as transfinite induction hypothesis $\forall \eta < \xi. f(\eta) = \eta$
 $\langle 1 \rangle 5$. $f(\xi) \subseteq \xi$
 $\langle 2 \rangle 1$. LET: $\eta \in f(\xi)$
 $\langle 2 \rangle 2$. PICK $\zeta \in \alpha$ such that $f(\zeta) = \eta$
 $\langle 2 \rangle 3$. $\zeta \in \xi$
PROOF: Since $f(\zeta) \in f(\xi)$ and f is a similarity.
 $\langle 2 \rangle 4$. $f(\zeta) = \zeta$
PROOF: $\langle 1 \rangle 4$
 $\langle 2 \rangle 5$. $\eta = \zeta$
PROOF: $\langle 2 \rangle 2, \langle 2 \rangle 4$
 $\langle 2 \rangle 6$. $\eta \in \xi$
PROOF: $\langle 2 \rangle 3, \langle 2 \rangle 5$
 $\langle 1 \rangle 6$. $\xi \subseteq f(\xi)$
 $\langle 2 \rangle 1$. LET: $\eta \in \xi$
 $\langle 2 \rangle 2$. $\eta = f(\eta) \in f(\xi)$
 $\langle 1 \rangle 7$. $f(\xi) = \xi$
 \square

Proposition 7.9. *Let α and β be ordinals. Then the following are equivalent.*

1. $\alpha \in \beta$
2. $\alpha \subsetneq \beta$
3. β is a continuation of α .

PROOF:

$\langle 1 \rangle 1$. $1 \Rightarrow 3$
PROOF: If $\alpha \in \beta$ then $\alpha = s(\alpha)$.
 $\langle 1 \rangle 2$. $3 \Rightarrow 2$
PROOF: Immediate from definitions.
 $\langle 1 \rangle 3$. $2 \Rightarrow 1$
 $\langle 2 \rangle 1$. LET: γ be the least element of β such that $\gamma \notin \alpha$
 $\langle 2 \rangle 2$. $\alpha \subseteq \gamma$
 $\langle 3 \rangle 1$. LET: $\eta \in \alpha$
 $\langle 3 \rangle 2$. $\eta \subseteq \alpha$
 $\langle 3 \rangle 3$. $\gamma \notin \eta$
 $\langle 3 \rangle 4$. $\eta \in \gamma$ or $\eta = \gamma$
 $\langle 3 \rangle 5$. $\eta \neq \gamma$
PROOF: Since $\eta \in \alpha$ and $\gamma \notin \alpha$.
 $\langle 3 \rangle 6$. $\eta \in \gamma$
 $\langle 2 \rangle 3$. $\gamma \subseteq \alpha$
PROOF: For all $\eta \in \gamma$ we have $\eta \in \alpha$ by leastness of γ .
 $\langle 2 \rangle 4$. $\gamma = \alpha$
 $\langle 2 \rangle 5$. $\alpha \in \beta$
 \square

Proposition 7.10. *For any ordinal numbers α and β , either $\alpha = \beta$, or $\alpha < \beta$, or $\beta < \alpha$.*

PROOF:

- ⟨1⟩1. Either $\alpha = \beta$, or α is similar to an initial segment of β , or β is similar to an initial segment of α .
- ⟨1⟩2. CASE: α is similar to an initial segment of β .
 - ⟨2⟩1. PICK $\eta \in \beta$ such that $\alpha \sim s(\eta)$
 - ⟨2⟩2. $\alpha \sim \eta$
 - ⟨2⟩3. $\alpha = \eta$
 - PROOF: Proposition 7.8.
 - ⟨2⟩4. $\alpha \in \beta$
- ⟨1⟩3. CASE: β is similar to an initial segment of α .
 PROOF: Then $\beta \in \alpha$ similarly.

□

Proposition 7.11. *Every set of ordinals is well ordered by $<$.*

PROOF:

- ⟨1⟩1. LET: E be a set of ordinals.
- ⟨1⟩2. LET: A be a nonempty subset of E .
- ⟨1⟩3. PICK $\alpha \in A$
- ⟨1⟩4. CASE: $\alpha \cap A = \emptyset$
 PROOF: Then α is least in A .
- ⟨1⟩5. CASE: $\alpha \cap A \neq \emptyset$
 PROOF: Then $\alpha \cap A$ has a least element, which is least in A .

□

Definition 7.12 (Limit Ordinal). A *limit ordinal* is an ordinal number that is not 0 and not α^+ for any ordinal α .

Proposition 7.13. *For any set E of ordinal numbers, $\bigcup E$ is an ordinal and is the supremum of E .*

PROOF: Proposition 5.37. □

Theorem 7.14 (Burali-Forti Paradox). *There is no set whose members are exactly the ordinal numbers.*

PROOF: For any set of ordinals E , we have $(\bigcup E)^+$ is an ordinal that is not in E . □

Theorem 7.15 (Counting Theorem). *Every well ordered set is similar to a unique ordinal.*

PROOF:

- ⟨1⟩1. LET: X be a well ordered set.
- ⟨1⟩2. There exists an ordinal α such that $X \cong \alpha$.
 - ⟨2⟩1. For all $a \in X$, there exists a unique ordinal α such that $s(a) \cong \alpha$
 - ⟨3⟩1. LET: $a \in X$
 - ⟨3⟩2. ASSUME: as transfinite induction hypothesis that, for all $b < a$, there exists a unique ordinal β such that $s(b) \cong \beta$

$\langle 3 \rangle 3$. LET: $\alpha = \{\beta : \beta \text{ is an ordinal} \wedge \exists b < a. s(b) \cong \beta\}$
 PROOF: This is a set by the Axiom of Substitution.
 $\langle 3 \rangle 4$. α is an ordinal
 $\langle 4 \rangle 1$. LET: $\gamma \in \beta \in \alpha$
 $\langle 4 \rangle 2$. PICK $b < a$ and $f : s(b) \cong \beta$
 $\langle 4 \rangle 3$. PICK $c < b$ such that $f(c) = \gamma$
 $\langle 4 \rangle 4$. $f \upharpoonright s(c) : s(c) \cong \gamma$
 $\langle 3 \rangle 5$. $s(a) \cong \alpha$
 PROOF: The function $f : s(a) \rightarrow \alpha$ defined by $f(b)$ is the ordinal such that $s(b) \cong f(b)$ is a similarity.
 $\langle 3 \rangle 6$. α is unique.
 PROOF: Proposition 7.8.
 $\langle 2 \rangle 2$. LET: $\alpha = \{\beta : \beta \text{ is an ordinal} \wedge \exists a \in X. s(a) \cong \beta\}$
 PROOF: This is a set by the Axiom of Substitution.
 $\langle 2 \rangle 3$. α is an ordinal.
 PROOF: Similar.
 $\langle 2 \rangle 4$. $X \cong \alpha$
 PROOF: Similar.
 $\langle 1 \rangle 3$. For any ordinals α and β , if $X \cong \alpha$ and $X \cong \beta$ then $\alpha = \beta$.
 PROOF: Proposition 7.8.
 \square

7.1 Order on the Natural Numbers

Proposition 7.16. *For natural numbers m, n and k , if $m < n$ then $m + k < n + k$.*

PROOF:
 $\langle 1 \rangle 1$. LET: $m, n \in \omega$
 $\langle 1 \rangle 2$. ASSUME: $m < n$
 $\langle 1 \rangle 3$. $m + 0 < n + 0$
 $\langle 1 \rangle 4$. $\forall k \in \omega. m + k < n + k \Rightarrow m + k^+ < n + k^+$
 PROOF: By Proposition 6.7.
 \square

Proposition 7.17. *For natural numbers m, n and k , if $m < n$ and $k \neq 0$ then $mk < nk$.*

PROOF:
 $\langle 1 \rangle 1$. LET: $m, n \in \omega$
 $\langle 1 \rangle 2$. ASSUME: $m < n$
 $\langle 1 \rangle 3$. $m1 < n1$
 $\langle 1 \rangle 4$. For all $k \in \omega$, if $k \neq 0$ and $mk < nk$ then $m(k + 1) < n(k + 1)$

PROOF:

$$\begin{aligned}
m(k+1) &= mk + m \\
&< mk + n && \text{(Proposition 7.16)} \\
&< nk + n && \text{(Proposition 7.16)} \\
&= n(k+1)
\end{aligned}$$

□

Proposition 7.18. *Let n be a natural number. Let X be a proper subset of n . Then there exists $m < n$ such that $X \sim m$.*

PROOF:

⟨1⟩1. LET: $P(n)$ be the property: for every proper subset $X \subsetneq n$, there exists $m < n$ such that $X \sim m$.

⟨1⟩2. $P(0)$

PROOF: Vacuous.

⟨1⟩3. $\forall n \in \omega. P(n) \Rightarrow P(n+1)$

⟨2⟩1. LET: $n \in \omega$

⟨2⟩2. ASSUME: $P(n)$

⟨2⟩3. LET: X be a proper subset of $n+1$

⟨2⟩4. CASE: $X - \{n\} = n$

PROOF: Then $X = n$ so $X \sim n < n+1$.

⟨2⟩5. CASE: $X - \{n\} \subsetneq n$

⟨3⟩1. PICK $m < n$ such that $X - \{n\} \sim m$

⟨3⟩2. $X \sim m$ or $X \sim m+1$

PROOF: If $n \in X$ then $X \sim m+1$. If $n \notin X$ then $X \sim m$.

□

Proposition 7.19. *For every natural number n , we have n is not equivalent to a proper subset of n .*

PROOF:

⟨1⟩1. LET: $P(n)$ be the property: every one-to-one function $n \rightarrow n$ is onto.

⟨1⟩2. $P(0)$

PROOF: The only function $0 \rightarrow 0$ is \emptyset .

⟨1⟩3. $\forall n \in \omega. P(n) \Rightarrow P(n+1)$

⟨2⟩1. LET: $n \in \omega$

⟨2⟩2. ASSUME: $P(n)$

⟨2⟩3. ASSUME: $f : n+1 \rightarrow n+1$ is one-to-one.

⟨2⟩4. LET: $g : n \rightarrow n$ be the function

$$g(k) = \begin{cases} f(k) & \text{if } f(k) < n \\ f(n) & \text{if } f(k) = n \end{cases}$$

PROOF: If $k < n$ and $f(k) = n$ then $f(n) < n$ since f is one-to-one.

⟨2⟩5. g is one-to-one.

⟨3⟩1. LET: $k, l < n$

⟨3⟩2. ASSUME: $g(k) = g(l)$

⟨3⟩3. CASE: $f(k) < n$ and $f(l) < n$

PROOF: Then $f(k) = g(k) = g(l) = f(l)$ so $k = l$ since f is one-to-one.

⟨3⟩4. CASE: $f(k) < n$ and $f(l) = n$
PROOF: Then $f(k) = g(k) = g(l) = f(n)$ contradicting the fact that f is one-to-one.

⟨3⟩5. CASE: $f(k) = n$ and $f(l) < n$
PROOF: Similar.

⟨3⟩6. CASE: $f(k) = n$ and $f(l) = n$
PROOF: Then $k = l$ since f is one-to-one.

⟨2⟩6. g maps n onto n .
PROOF: ⟨2⟩2

⟨2⟩7. f maps $n + 1$ onto $n + 1$.
⟨3⟩1. LET: $l < n + 1$
⟨3⟩2. CASE: $l < n$
⟨4⟩1. PICK $k < n$ such that $g(k) = l$
⟨4⟩2. $f(k) = l$ or $f(n) = l$
⟨3⟩3. CASE: $l = n$
⟨4⟩1. CASE: $f(n) = n$
PROOF: Then $l \in \text{ran } f$ as required.
⟨4⟩2. CASE: $f(n) < n$
⟨5⟩1. PICK $k < n$ such that $g(k) = f(n)$
⟨5⟩2. $f(k) = n$

□

Corollary 7.19.1. *Equivalent natural numbers are equal.*

Proposition 7.20. *The lexicographical order is a well ordering on $\omega \times \omega$.*

PROOF: Easy. □

7.2 Finite Sets

Definition 7.21 (Finite). A set is *finite* iff it is equivalent to a natural number; otherwise, it is *infinite*.

Proposition 7.22. *No finite set is equivalent to one of its proper subsets.*

PROOF: From Proposition 7.19. □

Proposition 7.23. *ω is infinite.*

PROOF: Since the function that maps n to $n + 1$ is a one-to-one correspondence between ω and $\omega - \{0\}$. □

Proposition 7.24. *Every subset of a finite set is finite.*

PROOF: Proposition 7.18. □

Definition 7.25 (Number of Elements). For any finite set E , the *number of elements* in E , $\sharp(E)$, is the unique natural number such that $E \sim \sharp(E)$.

Proposition 7.26. *Let E and F be finite sets. If $E \subseteq F$ then $\sharp(E) \leq \sharp(F)$.*

PROOF: Proposition 7.18. \square

Proposition 7.27. *Let E and F be disjoint finite sets. Then $E \cup F$ is finite and $\sharp(E \cup F) = \sharp(E) + \sharp(F)$.*

PROOF:

$\langle 1 \rangle 1$. LET: $P(n)$ be the statement: $n \in \omega$ and for any $m \in \omega$, if $E \sim m$, $F \sim n$ and $E \cap F = \emptyset$, then $E \cup F \sim m + n$

$\langle 1 \rangle 2$. $P(0)$

$\langle 2 \rangle 1$. LET: $m \in \omega$

$\langle 2 \rangle 2$. LET: $E \sim m$ and $F \sim 0$

$\langle 2 \rangle 3$. $F = \emptyset$

$\langle 2 \rangle 4$. $E \cup F = E \sim m = m + 0$

$\langle 1 \rangle 3$. $\forall n \in \omega. P(n) \Rightarrow P(n + 1)$

$\langle 2 \rangle 1$. LET: $n \in \omega$

$\langle 2 \rangle 2$. ASSUME: $P(n)$

$\langle 2 \rangle 3$. LET: $m \in \omega$

$\langle 2 \rangle 4$. LET: $E \sim m$ and $F \sim n + 1$

$\langle 2 \rangle 5$. ASSUME: $E \cap F = \emptyset$

$\langle 2 \rangle 6$. PICK $f \in F$

$\langle 2 \rangle 7$. $F - \{f\} \sim n$

$\langle 2 \rangle 8$. $E \cap (F - \{f\}) = \emptyset$

$\langle 2 \rangle 9$. $E \cup (F - \{f\}) \sim m + n$

PROOF: $\langle 2 \rangle 2$

$\langle 2 \rangle 10$. $E \cup F \sim m + n + 1$

\square

Corollary 7.27.1. *The union of two finite sets is finite.*

PROOF: Since, if E and F are finite, then $E \cup F = (E - F) \cup (E \cap F) \cup (F - E)$ and these are finite and disjoint. \square

Proposition 7.28. *If E and F are finite sets then $E \times F$ is finite and $\sharp(E \times F) = \sharp(E)\sharp(F)$.*

PROOF:

$\langle 1 \rangle 1$. LET: $P(n)$ be the statement: $n \in \omega$ and for all $m \in \omega$, if $E \sim m$ and $F \sim n$ then $E \times F \sim mn$

$\langle 1 \rangle 2$. $P(0)$

PROOF: If $F \sim 0$ then $F = \emptyset$ so $E \times F = \emptyset \sim 0$.

$\langle 1 \rangle 3$. $\forall n \in \omega. P(n) \Rightarrow P(n + 1)$

$\langle 2 \rangle 1$. LET: $n \in \omega$

$\langle 2 \rangle 2$. ASSUME: $P(n)$

$\langle 2 \rangle 3$. LET: $m \in \omega$

$\langle 2 \rangle 4$. ASSUME: $E \sim m$ and $F \sim n + 1$

$\langle 2 \rangle 5$. PICK $f \in F$

- ⟨2⟩6. $F - \{f\} \sim n$
- ⟨2⟩7. $E \times (F - \{f\}) \sim mn$
- ⟨2⟩8. $E \times F = (E \times (F - \{f\})) \cup (E \times \{f\})$
- ⟨2⟩9. $E \times \{f\} \sim m$
- ⟨2⟩10. $E \times F \sim mn + m$

PROOF: Proposition 7.27.

□

Proposition 7.29. *For any finite sets E and F , we have E^F is finite and $\sharp(E^F) = \sharp(E)^{\sharp(F)}$.*

PROOF:

- ⟨1⟩1. LET: $P(n)$ be the property: $n \in \omega$ and for all $m \in \omega$, if $E \sim m$ and $F \sim n$ then $E^F \sim m^n$
- ⟨1⟩2. $P(0)$
PROOF: Since $E^\emptyset = \{\emptyset\} \sim 1$
- ⟨1⟩3. $\forall n \in \omega. P(n) \Rightarrow P(n+1)$
 - ⟨2⟩1. LET: $n \in \omega$
 - ⟨2⟩2. ASSUME: $P(n)$
 - ⟨2⟩3. LET: $m \in \omega$
 - ⟨2⟩4. LET: $E \sim m$ and $F \sim n+1$
 - ⟨2⟩5. PICK $f \in F$
 - ⟨2⟩6. $F - \{f\} \sim n$
 - ⟨2⟩7. LET: $\phi : E^F \rightarrow E^{F-\{f\}} \times E$ be the function $\phi(g) = (g \upharpoonright (F - \{f\}), g(f))$
 - ⟨2⟩8. ϕ is a one-to-one correspondence
 - ⟨2⟩9. $\sharp(E^F) = m^{n+1}$

PROOF:

$$\begin{aligned}
 \sharp(E^F) &= \sharp(E^{F-\{f\}} \times E) \\
 &= \sharp(E^{F-\{f\}}) \sharp(E) && \text{(Proposition 7.28)} \\
 &= m^n m && (\langle 2 \rangle 2, \langle 2 \rangle 4) \\
 &= m^{n+1}
 \end{aligned}$$

□

Corollary 7.29.1. *If E is finite then $\mathcal{P}E$ is finite and $\sharp(\mathcal{P}E) = 2^{\sharp(E)}$.*

Proposition 7.30. *The union of a finite set of finite sets is finite.*

PROOF:

- ⟨1⟩1. LET: $P(n)$ be the property: for any set E , if $E \sim n$ and every element of E is finite, then $\bigcup E$ is finite.
- ⟨1⟩2. $P(0)$
PROOF: Since $\bigcup \emptyset = \emptyset$ is finite.
- ⟨1⟩3. $\forall n \in \omega. P(n) \Rightarrow P(n+1)$
 - ⟨2⟩1. LET: n be a natural number.
 - ⟨2⟩2. ASSUME: $P(n)$
 - ⟨2⟩3. LET: $E \sim n+1$

- ⟨2⟩4. PICK $X \in E$
- ⟨2⟩5. $E - \{X\} \sim n$
- ⟨2⟩6. $\bigcup(E - \{X\})$ is finite.
- PROOF: ⟨2⟩2
- ⟨2⟩7. $\bigcup E = \bigcup(E - \{X\}) \cup X$
- ⟨2⟩8. $\bigcup E$ is finite.
- PROOF: Corollary 7.27.1.

□

Proposition 7.31. *Every nonempty finite set of natural numbers has a greatest element.*

PROOF:

- ⟨1⟩1. LET: $P(n)$ be the property: for every $E \subseteq \mathbb{N}$, if $E \sim n$ then E has a greatest element.
- ⟨1⟩2. $P(1)$
- PROOF: Since k is the greatest element of $\{k\}$.
- ⟨1⟩3. $\forall n \geq 1. P(n) \Rightarrow P(n+1)$
- ⟨2⟩1. LET: $n \geq 1$
- ⟨2⟩2. ASSUME: $P(n)$
- ⟨2⟩3. ASSUME: $E \subseteq \omega$ and $E \sim n+1$
- ⟨2⟩4. PICK $k \in E$
- ⟨2⟩5. LET: l be the greatest element of $E - \{k\}$
- ⟨2⟩6. Either k or l is greatest in E .

□

Proposition 7.32. *Every infinite set has a subset equivalent to ω .*

PROOF:

- ⟨1⟩1. LET: X be an infinite set.
- ⟨1⟩2. PICK a choice function f for X .
- ⟨1⟩3. LET: \mathcal{C} be the set of all finite subsets of X .
- ⟨1⟩4. For all $A \in \mathcal{C}$ we have $X - A \in \text{dom } f$.
- PROOF: For all $A \in \mathcal{C}$ we have $X - A \neq \emptyset$.
- ⟨1⟩5. LET: $U : \omega \rightarrow \mathcal{C}$ be the function defined recursively by $U(0) = \emptyset$ and $U(n+1) = U(n) \cup \{f(X - U(n))\}$ for all $n \in \omega$.
- ⟨1⟩6. LET: $v : \omega \rightarrow X$ be the function $v(n) = f(X - U(n))$
- PROVE: v is one-to-one.
- ⟨1⟩7. $\forall n \in \omega. v(n) \notin U(n)$
- PROOF: Since $v(n) = f(X - U(n)) \in X - U(n)$.
- ⟨1⟩8. $\forall n \in \omega. v(n) \in U(n+1)$
- ⟨1⟩9. $\forall m, n \in \omega. n \leq m \Rightarrow U(n) \subseteq U(m)$
- PROOF: Since $U(n) \subseteq U(n+1)$ for all n .
- ⟨1⟩10. $\forall m, n \in \omega. n < m \Rightarrow v(n) \neq v(m)$
- PROOF: Since $v(n) \in U(m)$ and $v(m) \notin U(m)$.

□

Corollary 7.32.1. *A set is infinite if and only if it is equivalent to a proper subset.*

7.3 Ordinal Arithmetic

Definition 7.33 (Addition). Let I be a well ordered set and $(\alpha_i)_{i \in I}$ be a sequence of ordinals. Choose a well ordered set A_i such that $A_i \cong \alpha_i$ for each $i \in I$, and assume the sets A_i are pairwise disjoint. The *sum* $\sum_{i \in I} \alpha_i$ is the ordinal of the well ordered set $\bigcup_{i \in I} A_i$, where:

- for $x, y \in A_i$, we have $x <_{\bigcup_{i \in I} A_i} y$ if and only if $x <_{A_i} y$
- for $x \in A_i$ and $y \in A_j$ with $i \neq j$, we have $x <_{\bigcup_{i \in I} A_i} y$ iff $i <_I j$

We write $\alpha + \beta$ for $\sum_{i \in 2} \gamma_i$ where $\gamma_0 = \alpha$ and $\gamma_1 = \beta$.

Proposition 7.34.

$$\begin{aligned}\alpha + 0 &= \alpha \\ 0 + \alpha &= \alpha \\ \alpha + 1 &= \alpha^+ \\ \alpha + (\beta + \gamma) &= (\alpha + \beta) + \gamma\end{aligned}$$

PROOF: Easy. \square

Proposition 7.35. For any ordinals α and β , we have $\alpha < \beta$ if and only if there exists $\gamma \neq 0$ such that $\beta = \alpha + \gamma$.

PROOF: Easy. \square

Proposition 7.36.

$$1 + \omega = \omega$$

PROOF: Easy. \square

Definition 7.37 (Multiplication). Given ordinals α and β , the *product* $\alpha\beta$ is the ordinal of $\alpha \times \beta$ under the *reverse lexicographic order*: $(a, b) < (c, d)$ iff $b < d$ or $(b = d \text{ and } a < c)$.

Proposition 7.38.

$$\begin{aligned}\alpha 0 &= 0 \\ 0 \alpha &= 0 \\ \alpha 1 &= \alpha \\ 1 \alpha &= \alpha \\ \alpha(\beta \gamma) &= (\alpha \beta) \gamma \\ \alpha(\beta + \gamma) &= \alpha \beta + \alpha \gamma\end{aligned}$$

PROOF: Easy. \square

Proposition 7.39. For ordinals α and β , if $\alpha\beta = 0$ then $\alpha = 0$ or $\beta = 0$.

PROOF: Easy. \square

Example 7.40. The commutative law fails:

$$2\omega = \omega \neq \omega 2$$

PROOF: Easy. \square

Example 7.41. The right distributive law fails:

$$(1 + 1)\omega = \omega \neq 1\omega + 1\omega = \omega 2$$

Definition 7.42 (Exponentiation). Given ordinals α and β , define the ordinal α^β by

$$\begin{aligned} \alpha^0 &= 1 \\ \alpha^{\beta+1} &= \alpha^\beta \alpha \\ \alpha^\lambda &= \bigcup_{\beta < \lambda} \alpha^\beta \quad (\lambda \text{ a limit ordinal}) \end{aligned}$$

Proposition 7.43.

$$\begin{aligned} 0^\alpha &= 0 & (\alpha \geq 1) \\ 1^\gamma &= 1 \\ \alpha^{\beta+\gamma} &= \alpha^\beta \alpha^\gamma \\ \alpha^{\beta\gamma} &= (\alpha^\beta)^\gamma \end{aligned}$$

PROOF: Easy. \square

Example 7.44. $(\alpha\beta)^\gamma$ is different from $\alpha^\gamma\beta^\gamma$ in general:

$$(2 \cdot 2)^\omega = \omega \neq 2^\omega 2^\omega = \omega^2 .$$

7.4 Arithmetic on the Natural Numbers

Proposition 7.45. For all $m, n \in \omega$, we have

$$m + n = n + m .$$

PROOF:

$\langle 1 \rangle 1$. LET: $P(m)$ be the property $\forall n \in \omega. m + n = n + m$

$\langle 1 \rangle 2$. $P(0)$

$\langle 2 \rangle 1$. LET: $Q(n)$ be the property $0 + n = n + 0$

$\langle 2 \rangle 2$. $Q(0)$

PROOF: Trivial.

$\langle 2 \rangle 3$. $\forall n \in \omega. Q(n) \Rightarrow Q(n^+)$

$\langle 3 \rangle 1$. LET: $n \in \omega$

$\langle 3 \rangle 2$. ASSUME: $Q(n)$
 $\langle 3 \rangle 3$. $0 + n^+ = n^+ + 0$

PROOF:

$$\begin{aligned} 0 + n^+ &= (0 + n)^+ \\ &= (n + 0)^+ & (\langle 3 \rangle 2) \\ &= n^+ \\ &= n^+ + 0 \end{aligned}$$

$\langle 1 \rangle 3$. $\forall m \in \omega. P(m) \Rightarrow P(m^+)$

$\langle 2 \rangle 1$. LET: $m \in \omega$

$\langle 2 \rangle 2$. ASSUME: $P(m)$

$\langle 2 \rangle 3$. LET: $Q(n)$ be the property $m^+ + n = n + m^+$

$\langle 2 \rangle 4$. $Q(0)$

PROOF: $\langle 1 \rangle 2$

$\langle 2 \rangle 5$. $\forall n \in \omega. Q(n) \Rightarrow Q(n^+)$

$\langle 3 \rangle 1$. LET: $n \in \omega$

$\langle 3 \rangle 2$. ASSUME: $Q(n)$

$\langle 3 \rangle 3$. $Q(n^+)$

PROOF:

$$\begin{aligned} m^+ + n^+ &= (m^+ + n)^+ \\ &= (n + m^+)^+ & (\langle 3 \rangle 2) \\ &= (n + m)^{++} \\ &= (m + n)^{++} & (\langle 2 \rangle 2) \\ &= (m + n^+)^+ \\ &= (n^+ + m)^+ & (\langle 2 \rangle 2) \\ &= n^+ + m^+ \end{aligned}$$

□

Proposition 7.46. *For all $m, n \in \omega$, we have*

$$mn = nm .$$

PROOF:

$\langle 1 \rangle 1$. LET: $P(m)$ be the statement $\forall n \in \omega. mn = nm$

$\langle 1 \rangle 2$. $P(0)$

$\langle 2 \rangle 1$. LET: $Q(n)$ be the statement $0n = n0$

$\langle 2 \rangle 2$. $Q(0)$

PROOF: Trivial.

$\langle 2 \rangle 3$. $\forall n \in \omega. Q(n) \Rightarrow Q(n^+)$

$\langle 3 \rangle 1$. LET: $n \in \omega$

$\langle 3 \rangle 2$. ASSUME: $Q(n)$

$\langle 3 \rangle 3$. $Q(n^+)$

PROOF:

$$\begin{aligned}
 0n^+ &= 0n + 0 \\
 &= 0n \\
 &= n0 & (\langle 3 \rangle 2) \\
 &= 0 \\
 &= n^+0
 \end{aligned}$$

$\langle 1 \rangle 3. \forall m \in \omega. P(m) \Rightarrow P(m^+)$

$\langle 2 \rangle 1. \text{ LET: } m \in \omega$

$\langle 2 \rangle 2. \text{ ASSUME: } P(m)$

$\langle 2 \rangle 3. \text{ LET: } Q(n) \text{ be the statement } m^+n = nm^+$

$\langle 2 \rangle 4. Q(0)$

PROOF: $\langle 1 \rangle 2$

$\langle 2 \rangle 5. \forall n \in \omega. Q(n) \Rightarrow Q(n^+)$

$\langle 3 \rangle 1. \text{ LET: } n \in \omega$

$\langle 3 \rangle 2. \text{ ASSUME: } Q(n)$

$\langle 3 \rangle 3. Q(n^+)$

PROOF:

$$\begin{aligned}
 m^+n^+ &= m^+n + m^+ \\
 &= (m^+n + m)^+ \\
 &= (nm^+ + m)^+ & (\langle 3 \rangle 2) \\
 &= (nm + n + m)^+ \\
 &= (mn + m + n)^+ & (\langle 2 \rangle 2, \text{ Proposition 7.45}) \\
 &= (mn^+ + n)^+ \\
 &= (n^+m + n)^+ & (\langle 2 \rangle 2) \\
 &= n^+m + n^+ \\
 &= n^+m^+
 \end{aligned}$$

□

Chapter 8

Countable Sets

Definition 8.1 (Countable). A set A is *countable* or *denumerable* iff there exists an injective function $A \rightarrow \omega$.

Definition 8.2 (Countably Infinite). A set is *countably infinite* iff it is similar to ω .

Proposition 8.3. *Every subset of a countable set is countable.*

PROOF: Easy. \square

Proposition 8.4. *Let X be a set. If there exists a function from ω onto X , then X is countable.*

PROOF:

$\langle 1 \rangle 1$. LET: f be a function from ω onto X .

$\langle 1 \rangle 2$. Choose a function $g : X \rightarrow \omega$ such that, for all $x \in X$, we have $f(g(x)) = x$.

$\langle 1 \rangle 3$. g is one-to-one.

\square

Proposition 8.5. $\omega \times \omega$ is countable.

PROOF: The sequence

$(0, 0), (0, 1), (1, 0), (0, 2), (1, 1), (2, 0), \dots$

is an enumeration of $\omega \times \omega$. \square

Corollary 8.5.1. *A countable union of countable sets is countable.*

PROOF:

$\langle 1 \rangle 1$. LET: A be a countable set of countable sets.

$\langle 1 \rangle 2$. PICK a surjection $f : \omega \rightarrow A$

$\langle 1 \rangle 3$. For $n \in \omega$, PICK a surjection $g_n : \omega \rightarrow f(n)$

$\langle 1 \rangle 4$. PICK a surjection $h : \omega \rightarrow \omega \times \omega$

$\langle 1 \rangle 5$. $\lambda n \in \omega. g_{\pi_1(h(n))}(\pi_2(h(n)))$ is a surjection $\omega \rightarrow \bigcup A$

\square

Corollary 8.5.2. *The Cartesian product of two countable sets is countable.*

Corollary 8.5.3. *For any countable set A , the set of all finite subsets of A is countable.*

PROOF: Prove by induction on n that the set of all subsets of size n is countable. The set of all finite subsets is then the union of these. \square

Proposition 8.6. *$\mathcal{P}\omega$ is uncountable.*

PROOF: Cantor's Theorem. \square

Chapter 9

Cardinal Numbers

Definition 9.1 (Cardinal Number). A *cardinal number* or *initial ordinal* is an ordinal α such that, for all $\beta < \alpha$, we have $\beta \not\sim \alpha$.

Definition 9.2 (Cardinality). For any set X , the *cardinality* of X , $\text{card } X$, is the least ordinal that is equivalent to X .

Proposition 9.3. *Given sets X and Y , we have $X \sim Y$ if and only if $\text{card } X = \text{card } Y$.*

PROOF: Easy. \square

Proposition 9.4. *For sets X and Y , we have $\text{card } X \leq \text{card } Y$ if and only if there exists an injective function $X \rightarrow Y$.*

PROOF: Easy. \square

Proposition 9.5. *Every natural number is a cardinal. ω is a cardinal.*

PROOF: Easy. \square

Proposition 9.6. *Every infinite cardinal is a limit ordinal.*

PROOF: For α infinite we have $f : \alpha^+ \sim \alpha$ where $f(\alpha) = 0$ and $f(\beta) = \beta^+$ for all other β . \square

9.1 Cardinal Arithmetic

Definition 9.7 (Addition). Given a family of cardinal numbers $\{\kappa_i\}_{i \in I}$, let $\sum_{i \in I} \kappa_i$ be $\text{card} \bigcup_{i \in I} A_i$, where $\{A_i\}_{i \in I}$ is a pairwise disjoint family of sets with $\text{card } A_i = \kappa_i$ for all i .

We write $\kappa + \lambda$ for $\sum_{i \in 2} \kappa_i$ where $\kappa_0 = \kappa$ and $\kappa_1 = \lambda$.

Proposition 9.8.

$$\begin{aligned}\kappa + \lambda &= \lambda + \kappa \\ \kappa + (\lambda + \mu) &= (\kappa + \lambda) + \mu\end{aligned}$$

PROOF: Easy. \square

Proposition 9.9. *Cardinal addition agrees with ordinal addition on the natural numbers.*

PROOF: Easy induction. \square

Proposition 9.10. *If $\kappa \leq \kappa'$ then $\kappa + \lambda \leq \kappa' + \lambda$.*

PROOF: Easy. \square

Proposition 9.11. *If κ is an infinite cardinal number then $\kappa + \kappa = \kappa$.*

PROOF:

$\langle 1 \rangle 1$. LET: A be an infinite set.

PROVE: $A \times 2 \sim A$

$\langle 1 \rangle 2$. LET: \mathcal{F} be the set of all functions f such that there exists $X \subseteq A$ such that $f : X \times 2 \sim X$.

$\langle 1 \rangle 3$. \mathcal{F} is non-empty.

PROOF: Pick a subset $X \subseteq A$ such that $X \sim \omega$, and a bijection $X \times 2 \sim X$.

$\langle 1 \rangle 4$. \mathcal{F} is partially ordered by extension.

$\langle 1 \rangle 5$. Every chain in \mathcal{F} has an upper bound.

PROOF: If $\mathcal{C} \subseteq \mathcal{F}$ is a chain then $\bigcup \mathcal{C} \in \mathcal{F}$.

$\langle 1 \rangle 6$. PICK $f \in \mathcal{F}$ maximal.

$\langle 1 \rangle 7$. PICK $X \subseteq A$ such that $f : X \times 2 \sim X$

$\langle 1 \rangle 8$. $X - A$ is finite.

$\langle 2 \rangle 1$. ASSUME: for a contradiction $X - A$ is infinite.

$\langle 2 \rangle 2$. PICK $Y \subseteq X - A$ such that $Y \sim \omega$.

$\langle 2 \rangle 3$. PICK $g : Y \times 2 \sim Y$

$\langle 2 \rangle 4$. $f \cup g : (X \cup Y) \times 2 \sim X \cup Y$

$\langle 2 \rangle 5$. Q.E.D.

PROOF: This contradicts the maximality of f .

$\langle 1 \rangle 9$. $\text{card } A + \text{card } A = \text{card } A$

PROOF:

$$\begin{aligned}
 2 \text{ card } A &= 2(\text{card } X + \text{card}(A - X)) \\
 &= 2 \text{ card } X + 2 \text{ card}(A - X) \\
 &= \text{card } X + 2 \text{ card}(A - X) && (\langle 1 \rangle 7) \\
 &= \text{card } X && (\langle 1 \rangle 8) \\
 &= \text{card } X + \text{card}(A - X) && (\langle 1 \rangle 8) \\
 &= \text{card } A
 \end{aligned}$$

\square

Corollary 9.11.1. *For any cardinals κ and λ that are not both finite, we have*

$$\kappa + \lambda = \max(\kappa, \lambda) .$$

Definition 9.12 (Multiplication). Given a family of cardinal numbers $\{\kappa_i\}_{i \in I}$, let $\prod_{i \in I} \kappa_i = \text{card} \times_{i \in I} \kappa_i$.

We write $\kappa \lambda$ for $\prod_{i \in 2} \kappa_i$ where $\kappa_0 = \kappa$ and $\kappa_1 = \lambda$.

Proposition 9.13.

$$\begin{aligned}\kappa\lambda &= \lambda\kappa \\ \kappa(\lambda\mu) &= (\kappa\lambda)\mu \\ \kappa(\lambda + \mu) &= \kappa\lambda + \kappa\mu\end{aligned}$$

Proposition 9.14. *Cardinal multiplication agrees with ordinal multiplication on the natural numbers.*

PROOF: Easy induction. \square

Proposition 9.15. *If $\kappa \leq \kappa'$ then $\kappa\lambda \leq \kappa'\lambda$.*

PROOF: Easy. \square

Proposition 9.16. *Let $\{\kappa_i\}_{i \in I}$ and $\{\lambda_i\}_{i \in I}$ be families of cardinal numbers with the same index set. If $\kappa_i < \lambda_i$ for all i , then $\sum_{i \in I} \kappa_i < \prod_{i \in I} \lambda_i$.*

PROOF:

$\langle 1 \rangle 1$. Choose a one-to-one function $f_i : \kappa_i \rightarrow \lambda_i$ for each $i \in I$

$\langle 1 \rangle 2$. $\sum_{i \in I} \kappa_i \leq \prod_{i \in I} \lambda_i$

PROOF: Define $g : \sum_{i \in I} \kappa_i \rightarrow \prod_{i \in I} \lambda_i$ by

$$g(i, \eta)(j) = \begin{cases} f_i(\eta) & \text{if } i = j \\ 0 & \text{if } i \neq j \end{cases}$$

$\langle 1 \rangle 3$. There is no surjective function $\sum_{i \in I} \kappa_i \rightarrow \prod_{i \in I} \lambda_i$

$\langle 2 \rangle 1$. LET: $h : \sum_i \kappa_i \rightarrow \prod_i \lambda_i$

$\langle 2 \rangle 2$. Choose $t(i) < \lambda_i$ for each $i \in I$ such that, for all $\eta < \kappa_i$, we have $t(i) \neq h(i, \eta)(i)$.

PROOF: Since the function that maps η to $h(i, \eta)(i)$ cannot be surjective $\kappa_i \rightarrow \lambda_i$.

$\langle 2 \rangle 3$. For all $i \in I$ and $\eta < \kappa_i$ we have $h \neq t(i, \eta)$.

\square

Proposition 9.17. *If κ is an infinite cardinal then $\kappa\kappa = \kappa$.*

PROOF:

$\langle 1 \rangle 1$. LET: A be an infinite set.

$\langle 1 \rangle 2$. LET: \mathcal{F} be the set of all functions f such that there exists $X \subseteq A$ such that $f : X \times X \sim X$

$\langle 1 \rangle 3$. \mathcal{F} is nonempty.

PROOF: Pick a countably infinite $X \subseteq A$. Then $X \times X \sim X$.

$\langle 1 \rangle 4$. \mathcal{F} is partially ordered by extension.

$\langle 1 \rangle 5$. Every chain in \mathcal{F} has an upper bound.

$\langle 1 \rangle 6$. PICK $f \in \mathcal{F}$ maximal.

$\langle 1 \rangle 7$. PICK $X \subseteq A$ such that $f : X \times X \sim X$.

$\langle 1 \rangle 8$. $\text{card } X = \text{card } A$

$\langle 2 \rangle 1$. ASSUME: for a contradiction $\text{card } X < \text{card } A$

$\langle 2 \rangle 2$. $\text{card } A = \text{card}(A - X)$

PROOF: Corollary 9.11.1.

$\langle 2 \rangle 3$. $\text{card } X < \text{card}(A - X)$

$\langle 2 \rangle 4$. PICK $Y \subseteq A - X$ such that $Y \sim X$

$\langle 2 \rangle 5$. PICK $g : (X \times Y) \cup (Y \times X) \cup (Y \times Y) \sim Y$

PROOF:

$$(X \times Y) \cup (Y \times X) \cup (Y \times Y) \sim 3 \times X \times X \quad (\langle 2 \rangle 4)$$

$$\sim 3 \times X \quad (\langle 1 \rangle 7)$$

$$\sim X \quad (\text{Corollary 9.11.1})$$

$$\sim Y \quad (\langle 2 \rangle 4)$$

$\langle 2 \rangle 6$. $f \cup g : (X \cup Y) \times (X \cup Y) \sim X \cup Y$

$\langle 2 \rangle 7$. Q.E.D.

PROOF: This contradicts the maximality of f .

□

Corollary 9.17.1. *If κ and λ are non-zero cardinals that are not both finite, then*

$$\kappa\lambda = \max(\kappa, \lambda) \text{ .}$$

Definition 9.18 (Exponentiation). Given cardinal numbers κ and λ , let κ^λ be the cardinality of the set of all functions $\lambda \rightarrow \kappa$.

Proposition 9.19.

$$\kappa^{\lambda+\mu} = \kappa^\lambda \kappa^\mu$$

$$(\kappa\lambda)^\mu = \kappa^\mu \lambda^\mu$$

$$\kappa^{\lambda\mu} = (\kappa^\lambda)^\mu$$

PROOF: Easy. □

Proposition 9.20. *Cardinal exponentiation and ordinal exponentiation agree on the natural numbers.*

PROOF: Easy. □

Proposition 9.21.

$$\text{card } \mathcal{P}X = 2^{\text{card } X}$$

PROOF: Define $\chi : \mathcal{P}X \sim 2^X$ to be the function that maps S to the function $\chi_S : X \rightarrow 2$ where $\chi_S(x) = 1$ if $x \in S$ and $\chi_S(x) = 0$ if $x \notin S$. □

Proposition 9.22. *For any infinite cardinal κ we have $\kappa < 2^\kappa$.*

PROOF: Proposition 9.16. □

Proposition 9.23. *If $\kappa \leq \lambda$ then $\kappa^\mu \leq \lambda^\mu$.*

PROOF: Easy. □

9.2 Alephs

Definition 9.24 (Aleph). Define the cardinal \aleph_α for every ordinal α as follows: \aleph_α is the least infinite cardinal greater than \aleph_β for all $\beta < \alpha$.

Proposition 9.25.

$$\aleph_0 = \omega$$

PROOF: Easy. \square

Definition 9.26 (Continuum Hypothesis). The *continuum hypothesis* is the statement $\aleph_1 = 2^{\aleph_0}$.

Definition 9.27 (Generalized Continuum Hypothesis). The *generalized continuum hypothesis* is the statement: for every ordinal α we have $\aleph_{\alpha+1} = 2^{\aleph_\alpha}$.

Chapter 10

Field Theory

Definition 10.1 (Field). A *field* is a triple $(K, +, \cdot)$ such that K is a set, $+$ and \cdot are functions $K^2 \rightarrow K$, and:

1. $\forall x, y, z \in K. x + (y + z) = (x + y) + z$
2. $\forall x, y, z \in K. x(yz) = (xy)z$
3. $\forall x, y \in K. x + y = y + x$
4. $\forall x, y \in K. xy = yx$
5. There exists a unique $0 \in K$ such that $\forall x \in K. x + 0 = x$
6. There exists a unique $1 \in K$ such that $\forall x \in K. x1 = x$
7. $0 \neq 1$
8. $\forall x \in K. \exists!(-x) \in K. x + (-x) = 0$
9. For all $x \in K$, if $x \neq 0$ then $\exists!(1/x) \in K. x \cdot 1/x = 1$
10. $\forall x, y, z \in K. x(y + z) = xy + xz$

Definition 10.2 (Subtraction). In any field K , define *subtraction* $- : K^2 \rightarrow K$ by $x - y = x + (-y)$.

Definition 10.3 (Ordered Field). An *ordered field* is a quadruple $(K, +, \cdot, \leq)$ such that $(K, +, \cdot)$ is a field and \leq is a linear order on K , and:

1. For all $x, y, z \in K$, if $x < y$ then $x + z < y + z$
2. For all $x, y, z \in K$, if $x < y$ and $0 < z$ then $xz < yz$

Chapter 11

Real Numbers

11.1 Axioms for Real Numbers

Let there be a set \mathbb{R} , whose elements are called *real numbers*.

Let there be two functions $+, \cdot : \mathbb{R}^2 \rightarrow \mathbb{R}$.

Let there be a relation $< \subseteq \mathbb{R}^2$.

Axiom 11.1 (Associativity of Addition).

$$\forall x, y, z \in \mathbb{R}. x + (y + z) = (x + y) + z$$

Axiom 11.2 (Associativity of Multiplication).

$$\forall x, y, z \in \mathbb{R}. x(yz) = (xy)z$$

Axiom 11.3 (Commutativity of Addition).

$$\forall x, y \in \mathbb{R}. x + y = y + x$$

Axiom 11.4 (Commutativity of Multiplication).

$$\forall x, y \in \mathbb{R}. xy = yx$$

Axiom 11.5 (Identity for Addition). *There exists a unique $z \in \mathbb{R}$ such that $\forall x \in \mathbb{R}. x + z = x$.*

Definition 11.6 (Zero). The real number *zero*, 0, is the unique real number such that $\forall x \in \mathbb{R}. x + 0 = x$.

Axiom 11.7 (Identity for Multiplication). *There exists a unique $i \in \mathbb{R}$ such that $\forall x \in \mathbb{R}. xi = x$. Further, we have $i \neq 0$.*

Definition 11.8 (One). The real number *one*, 1, is the unique real number such that $\forall x \in \mathbb{R}. x1 = x$.

Axiom 11.9 (Additive Inverses). *For all $x \in \mathbb{R}$, there exists a unique $y \in \mathbb{R}$ such that $x + y = 0$.*

Axiom 11.10 (Multiplicative Inverses). *For all $x \in \mathbb{R}$, if $x \neq 0$ then there exists a unique $y \in \mathbb{R}$ such that $xy = 1$.*

Axiom 11.11 (Distributive Law).

$$\forall x, y, z \in \mathbb{R}. x(y + z) = xy + xz$$

Axiom 11.12 (Monotonicity of Addition). *For all $x, y, z \in \mathbb{R}$, if $x < y$ then $x + z < y + z$.*

Axiom 11.13 (Monotonicity of Multiplication). *For all $x, y, z \in \mathbb{R}$, if $x < y$ and $0 < z$ then $xz < yz$.*

Axiom 11.14 (Least Upper Bound Property). *The relation $<$ is a strict linear order on \mathbb{R} with the least upper bound property.*

11.2 Consequences of the Axioms

11.2.1 Negation

Definition 11.15 (Negation). For any real number x , the *negation* of x , $-x$, is the unique real number such that $x + (-x) = 0$.

Theorem 11.16. *For any real numbers x and y , if $x + y = x$ then $y = 0$.*

PROOF:

$\langle 1 \rangle 1.$ LET: $x, y \in \mathbb{R}$

$\langle 1 \rangle 2.$ ASSUME: $x + y = x$

$\langle 1 \rangle 3.$ $y = 0$

PROOF:

$$\begin{aligned} y &= y + 0 && \text{(Definition of zero)} \\ &= y + (x + (-x)) && \text{(Definition of } -x) \\ &= (y + x) + (-x) && \text{(Associativity of Addition)} \\ &= (x + y) + (-x) && \text{(Commutativity of Addition)} \\ &= x + (-x) && (\langle 1 \rangle 2) \\ &= 0 && \text{(Definition of } -x) \end{aligned}$$

□

Theorem 11.17.

$$\forall x \in \mathbb{R}. 0x = 0$$

PROOF:

$\langle 1 \rangle 1.$ LET: $x \in \mathbb{R}$

$\langle 1 \rangle 2.$ $xx + 0x = xx$

PROOF:

$$\begin{aligned} xx + 0x &= (x + 0)x && \text{(Distributive Law)} \\ &= xx && \text{(Definition of 0)} \end{aligned}$$

$\langle 1 \rangle 3. 0x = 0$

PROOF: Theorem 11.16, $\langle 1 \rangle 2$.

□

Theorem 11.18.

$$-0 = 0$$

PROOF: Since $0 + 0 = 0$. □

Theorem 11.19.

$$\forall x \in \mathbb{R}. -(-x) = x$$

PROOF: Since $-x + x = 0$. □

Theorem 11.20.

$$\forall x, y \in \mathbb{R}. x(-y) = -(xy)$$

PROOF:

$$\begin{aligned} x(-y) + xy &= x((-y) + y) && \text{(Distributive Law)} \\ &= x0 && \text{(Definition of } -y) \\ &= 0 && \text{(Theorem 11.17)} \end{aligned} \quad \square$$

Theorem 11.21.

$$\forall x \in \mathbb{R}. (-1)x = -x$$

PROOF:

$$\begin{aligned} (-1)x &= -(1 \cdot x) && \text{(Theorem 11.20)} \\ &= -x && \text{(Definition of 1)} \end{aligned} \quad \square$$

11.2.2 Subtraction

Theorem 11.22.

$$\forall x, y, z \in \mathbb{R}. x(y - z) = xy - xz$$

PROOF:

$$\begin{aligned} x(y - z) &= x(y + (-z)) && \text{(Definition of subtraction)} \\ &= xy + x(-z) && \text{(Distributive Law)} \\ &= xy + (-(xz)) && \text{(Theorem 11.20)} \\ &= xy - xz && \text{(Definition of subtraction)} \end{aligned} \quad \square$$

Theorem 11.23.

$$\forall x, y \in \mathbb{R}. -(x + y) = -x - y$$

PROOF:

$$\begin{aligned}
-(x + y) &= (-1)(x + y) && \text{(Theorem 11.21)} \\
&= (-1)x + (-1)y && \text{(Distributive Law)} \\
&= -x + (-y) && \text{(Theorem 11.21)} \\
&= -x - y && \text{(Definition of subtraction)} \quad \square
\end{aligned}$$

Theorem 11.24.

$$\forall x, y \in \mathbb{R}. -(x - y) = -x + y$$

PROOF:

$$\begin{aligned}
-(x - y) &= -(x + (-y)) && \text{(Definition of subtraction)} \\
&= -x - (-y) && \text{(Theorem 11.23)} \\
&= -x + (-(-y)) && \text{(Definition of subtraction)} \\
&= -x + y && \text{(Theorem 11.19)} \quad \square
\end{aligned}$$

Definition 11.25 (Reciprocal). Given $x \in \mathbb{R}$ with $x \neq 0$, the *reciprocal* of x , $1/x$, is the unique real number such that $x \cdot 1/x = 1$.

Theorem 11.26. For any real numbers x and y , if $x \neq 0$ and $xy = x$ then $y = 1$.

PROOF:

- $\langle 1 \rangle 1$. LET: $x, y \in \mathbb{R}$
- $\langle 1 \rangle 2$. ASSUME: $x \neq 0$
- $\langle 1 \rangle 3$. ASSUME: $xy = x$
- $\langle 1 \rangle 4$. $y = 1$

PROOF:

$$\begin{aligned}
y &= y1 && \text{(Definition of 1)} \\
&= y(x \cdot 1/x) && \text{(Definition of } 1/x, \langle 1 \rangle 2) \\
&= (yx)1/x && \text{(Associativity of Multiplication)} \\
&= (xy)1/x && \text{(Commutativity of Multiplication)} \\
&= x \cdot 1/x && (\langle 1 \rangle 3) \\
&= 1 && \text{(Definition of } 1/x, \langle 1 \rangle 2)
\end{aligned}$$

\square

Definition 11.27 (Quotient). Given real numbers x and y with $y \neq 0$, the *quotient* x/y is defined by

$$x/y = x \cdot 1/y .$$

Theorem 11.28. For any real number x , if $x \neq 0$ then $x/x = 1$.

PROOF: Immediate from definitions. \square

Theorem 11.29.

$$\forall x \in \mathbb{R}. x/1 = x$$

PROOF:

$\langle 1 \rangle 1$. LET: $x \in \mathbb{R}$

$\langle 1 \rangle 2$. $1/1 = 1$

PROOF: Since $1 \cdot 1 = 1$.

$\langle 1 \rangle 3$. $x/1 = x$

PROOF: Since $x/1 = x \cdot 1/1 = x \cdot 1 = x$.

□

Theorem 11.30. For any real numbers x and y , if $x \neq 0$ and $y \neq 0$ then $xy \neq 0$.

PROOF:

$\langle 1 \rangle 1$. LET: $x, y \in \mathbb{R}$

$\langle 1 \rangle 2$. ASSUME: $xy = 0$ and $x \neq 0$

PROVE: $y = 0$

$\langle 1 \rangle 3$. $y = 0$

PROOF:

$$\begin{aligned} y &= 1y && \text{(Definition of 1)} \\ &= (1/x)xy && \text{(Definition of } 1/x, \langle 1 \rangle 2) \\ &= (1/x)0 && (\langle 1 \rangle 2) \\ &= 0 && \text{(Theorem 11.17)} \end{aligned}$$

□

Theorem 11.31. For any real numbers y and z , if $y \neq 0$ and $z \neq 0$ then $(1/y)(1/z) = 1/(yz)$.

PROOF: Since $yz(1/y)(1/z) = 1 \cdot 1 = 1$. □

Corollary 11.31.1. For any real numbers x, y, z, w with $y \neq 0 \neq w$, we have $(x/y)(z/w) = (xz)/(yw)$.

Theorem 11.32. For any real numbers x, y, z, w with $y \neq 0 \neq w$, we have

$$\frac{x}{y} + \frac{z}{w} = \frac{xw + yz}{yw}$$

PROOF:

$$\begin{aligned} yw \left(\frac{x}{y} + \frac{z}{w} \right) &= yw \frac{x}{y} + yw \frac{z}{w} \\ &= wx + yz \end{aligned} \quad \square$$

Theorem 11.33. For any real number x , if $x \neq 0$ then $1/x \neq 0$.

PROOF: Since $x \cdot 1/x = 1 \neq 0$. □

Theorem 11.34. For any real numbers w, z , if $w \neq 0 \neq z$ then $1/(w/z) = z/w$.

PROOF: Since $(z/w)(w/z) = (wz)/(wz) = 1$. □

Theorem 11.35. For any real numbers a, x and y , if $y \neq 0$ then $(ax)/y = a(x/y)$

PROOF: Since $ya(x/y) = ax$. \square

Theorem 11.36. For any real numbers x and y , if $y \neq 0$ then $(-x)/y = x/(-y) = -(x/y)$.

PROOF:

$\langle 1 \rangle 1$. $(-x)/y = -(x/y)$

PROOF: Take $a = -1$ in Theorem 11.35.

$\langle 1 \rangle 2$. $x/(-y) = -(x/y)$

PROOF: Since $(-y)(-(x/y)) = y(x/y) = x$.

\square

Theorem 11.37. For any real numbers x, y, z and w , if $x > y$ and $w > z$ then $x + w > y + z$.

PROOF: We have $y + z < x + z < x + w$ by Monotonicity of Addition twice. \square

Corollary 11.37.1. For any real numbers x and y , if $x > 0$ and $y > 0$ then $x + y > 0$.

Theorem 11.38. For any real numbers x and y , if $x > 0$ and $y > 0$ then $xy > 0$.

PROOF:

$$\begin{aligned} xy &> 0y && \text{(Monotonicity of Multiplication)} \\ &= 0 && \text{(Theorem 11.17)} \end{aligned} \quad \square$$

Theorem 11.39. For any real number x , we have $x > 0$ iff $-x < 0$.

PROOF:

$\langle 1 \rangle 1$. If $0 < x$ then $-x < 0$

PROOF: By Monotonicity of Addition adding $-x$ to both sides.

$\langle 1 \rangle 2$. If $-x < 0$ then $0 < x$

PROOF: By Monotonicity of Addition adding x to both sides.

\square

Theorem 11.40. For any real numbers x and y , we have $x > y$ iff $-x < -y$.

PROOF:

$\langle 1 \rangle 1$. If $y < x$ then $-x < -y$.

PROOF: By Monotonicity of Addition adding $-x - y$ to both sides.

$\langle 1 \rangle 2$. If $-x < -y$ then $y < x$.

PROOF: By Monotonicity of Addition adding $x + y$ to both sides.

\square

Theorem 11.41. For any real numbers x, y and z , if $x > y$ and $z < 0$ then $xz < yz$.

PROOF:

$\langle 1 \rangle 1$. LET: x, y and z be real numbers.

$\langle 1 \rangle 2$. ASSUME: $x > y$

$\langle 1 \rangle 3$. ASSUME: $z < 0$

$\langle 1 \rangle 4$. $-z > 0$

PROOF: Theorem 11.39, $\langle 1 \rangle 3$.

$\langle 1 \rangle 5$. $x(-z) > y(-z)$

PROOF: $\langle 1 \rangle 2$, $\langle 1 \rangle 4$, Monotonicity of Multiplication.

$\langle 1 \rangle 6$. $-(xz) > -(yz)$

PROOF: Theorem 11.20, $\langle 1 \rangle 5$.

$\langle 1 \rangle 7$. $xz < yz$

PROOF: Theorem 11.39, $\langle 1 \rangle 6$.

□

Theorem 11.42. *For any real number x , if $x \neq 0$ then $xx > 0$.*

PROOF:

$\langle 1 \rangle 1$. If $x > 0$ then $xx > 0$

PROOF: By Monotonicity of Multiplication.

$\langle 1 \rangle 2$. If $x < 0$ then $xx > 0$

PROOF: Theorem 11.41.

□

Theorem 11.43.

$$0 < 1$$

PROOF: By Theorem 11.42 since $1 = 1 \cdot 1$. □

Definition 11.44 (Positive). A real number x is *positive* iff $x > 0$.

We write \mathbb{R}_+ for the set of positive reals.

Theorem 11.45. *For any real numbers x and y , we have xy is positive if and only if x and y are both positive or both negative.*

PROOF: By the Monotonicity of Multiplication and Theorem 11.41. □

Corollary 11.45.1. *For any real number x , if $x > 0$ then $1/x > 0$.*

PROOF: Since $x \cdot 1/x = 1$ is positive. □

Theorem 11.46. *For any real numbers x and y , if $x > y > 0$ then $1/x < 1/y$.*

PROOF: If $1/y \leq 1/x$ then $1 < 1$ by Monotonicity of Multiplication. □

Theorem 11.47. *For any real numbers x and y , if $x < y$ then $x < (x+y)/2 < y$.*

PROOF: We have $2x < x+y$ and $x+y < 2y$ by Monotonicity of Addition, hence $x < (x+y)/2 < y$ by Monotonicity of Multiplication since $1/2 > 0$. □

Corollary 11.47.1. \mathbb{R} is a linear continuum.

Definition 11.48 (Negative). A real number x is *negative* iff $x < 0$.

We write $\overline{\mathbb{R}_+}$ for the set of nonnegative reals.

Chapter 12

Integers

Definition 12.1 (Inductive). A set of real numbers A is *inductive* iff $1 \in A$ and $\forall x \in A. x + 1 \in A$.

Definition 12.2 (Positive Integer). The set \mathbb{Z}_+ of *positive integers* is the intersection of the set of inductive sets.

Proposition 12.3. *Every positive integer is positive.*

PROOF: The set of positive reals is inductive. \square

Proposition 12.4. *1 is the least element of \mathbb{Z}_+ .*

PROOF: Since $\{x \in \mathbb{R} : x \geq 1\}$ is inductive. \square

Proposition 12.5. *\mathbb{Z}_+ is inductive.*

PROOF: 1 is an element of every inductive set, and for all $x \in \mathbb{R}$, if x is an element of every inductive set then so is $x + 1$. \square

Theorem 12.6 (Principle of Induction). *If A is an inductive set of positive integers then $A = \mathbb{Z}_+$.*

PROOF: Immediate from definitions. \square

Definition 12.7 (Integer). The set \mathbb{Z} of *integers* is

$$\mathbb{Z} = \mathbb{Z}_+ \cup \{0\} \cup \{-x : x \in \mathbb{Z}_+\} .$$

Proposition 12.8. *The sum, difference and product of two integers is an integer.*

PROOF: Easy. \square

Example 12.9. $1/2$ is not an integer.

Definition 12.10 (Rational Number). The set \mathbb{Q} of *rational numbers* is the set of all real numbers that are the quotient of two integers.