Summary of Halmos' Naive Set Theory

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Contents

Chapter 1

Primitive Terms and Axioms

Let there be *sets*. We assume that everything is a set.

Let there be a binary relation of *membership*, $\in$. If $x \in A$ we say that x *belongs* to A, x is an *element* of A, or x is *contained* in A. If this does not hold we write $x \notin A$.

Axiom 1.1 (Axiom of Extensionality). *Two sets are equal if and only if they have the same elements.*

Axiom 1.2 (Axiom of Comprehension, Aussonderungsaxiom). *To every set A and to every condition $S(x)$ there corresponds a set B whose elements are exactly those elements x of A for which $S(x)$ holds.*

Axiom 1.3 (Axiom of Pairing). *For any two sets, there exists a set that they both belong to.*

Axiom 1.4 (Union Axiom). *For every set A, there exists a set that contains all the elements that belong to at least one element of A.*

Definition 1.5 (Subset). Let A and B be sets. We say that A is a *subset* of B, or B *includes* A, and write $A \subseteq B$ or $B \supseteq A$, iff every element of A is an element of B.

Axiom 1.6 (Power Set Axiom). *For any set A, there exists a set that contains all the subsets of A.*

Axiom 1.7 (Axiom of Infinity). *There exists a set I such that:*

- *I has an element that has no elements*

- *for all $x \in I$, there exists $y \in I$ such that the elements of y are exactly x and the elements of x.*

Chapter 2

Basic Properties and Operations on Sets

2.1 The Subset Relation

Theorem 2.1. *For any set A, we have $A \subseteq A$.*

PROOF: Every element of A is an element of A. □

Theorem 2.2. *For any sets A, B and C, if $A \subseteq B$ and $B \subseteq C$ then $A \subseteq C$.*

PROOF: If every element of A is an element of B, and every element of B is an element of C, then every element of A is an element of C. □

Theorem 2.3. *For any sets A and B, if $A \subseteq B$ and $B \subseteq A$ then $A = B$.*

PROOF: If every element of A is an element of B, and every element of B is an element of A, then A and B have the same elements, and therefore are equal by the Axiom of Extensionality. □

Definition 2.4 (Proper Subset)**.** Let A and B be sets. We say that A is a *proper* subset of B, or B *properly* includes A, and write $A \subsetneq B$ or $B \supsetneq A$, iff $A \subseteq B$ and $A \neq B$.

2.2 Comprehension Notation

Definition 2.5. Given a set A and a condition $S(x)$, we write $\{x \in A : S(x)\}$ for the set whose elements are exactly those elements x of A for which $S(x)$ holds.

PROOF: This exists by the Axiom of Comprehension and is unique by the Axiom of Extensionality. □

Theorem 2.6. *There is no set that contains every set.*

PROOF:

⟨1⟩1. LET: A be a set.

PROVE: There exists a set B such that $B \notin A$.

⟨1⟩2. LET: $B = \{x \in A : x \notin x\}$

⟨1⟩3. If $B \in A$ then we have $B \in B$ if and only if $B \notin B$.

⟨1⟩4. $B \notin A$

□

2.3 The Empty Set

Theorem 2.7. *There exists a set with no elements.*

PROOF: Immediate from the Axiom of Infinity. □

Definition 2.8 (Empty Set)**.** The *empty set* $\varnothing$ is the set with no elements.

Theorem 2.9. *For any set A we have $\varnothing \subset A$.*

PROOF: Vacuous. □

2.4 Unordered Pairs

Definition 2.10 ((Unordered) Pair)**.** For any sets a and b, the *(unordered) pair* $\{a, b\}$ is the set whose elements are just a and b.

PROOF: This exists by the Axioms of Pairing and Comprehension, and is unique by the Axiom of Extensionality. □

Definition 2.11 (Singleton)**.** For any set a, the *singleton* $\{a\}$ is defined to be $\{a, a\}$.

2.5 Unions

Definition 2.12 (Union)**.** For any set $\mathcal{C}$, the *union* of $\mathcal{C}$, $\bigcup \mathcal{C}$, is the set whose elements are the elements of the elements of $\mathcal{C}$.

We write $\bigcup_{X \in \mathcal{A}} t[X]$ for $\bigcup \{t[X] \mid X \in \mathcal{A}\}$.

PROOF: This exists by the Union Axiom and Comprehension Axiom, and is unique by the Axiom of Extensionality. □

Proposition 2.13.
$$\bigcup \varnothing = \varnothing$$

PROOF: There is no set that is an element of an element of $\varnothing$. □

Proposition 2.14. *For any set A, we have $\bigcup \{A\} = A$.*

PROOF: For any x, we have x is an element of an element of $\{A\}$ if and only if x is an element of A. □

Definition 2.15. We write $A \cup B$ for $\bigcup\{A, B\}$.

Proposition 2.16. *For any set A, we have $A \cup \varnothing = A$.*

PROOF: $x \in A \cup \varnothing$ iff $x \in A$ or $x \in \varnothing$, iff $x \in A$. □

Proposition 2.17 (Idempotence). *For any set A, we have $A \cup A = A$.*

PROOF: $x \in A$ or $x \in A$ is equivalent to $x \in A$. □

Proposition 2.18. *For any sets A and B, we have $A \subseteq B$ if and only if $A \cup B = B$.*

PROOF: For any x, the statement "if $x \in A$ then $x \in B$" is equivalent to "$x \in A$ or $x \in B$ if and only if $x \in B$". □

Proposition 2.19. *For any sets a and b, we have $\{a\} \cup \{b\} = \{a, b\}$.*

PROOF: Immediate from definitions. □

2.6 Intersections

Definition 2.20 (Intersection). *For any sets A and B, the *intersection* $A \cap B$ is defined to be $\{x \in A : x \in B\}$.*

Proposition 2.21. *For any set A, we have $A \cap \varnothing = \varnothing$.*

PROOF: There is no x such that $x \in A$ and $x \in \varnothing$. □

Proposition 2.22. *For any set A, we have*

$$A \cap A = A \ .$$

PROOF: We have $x \in A$ and $x \in A$ if and only if $x \in A$. □

Proposition 2.23. *For any sets A and B, we have $A \subseteq B$ if and only if $A \cap B = A$.*

PROOF: For any x, the statement "if $x \in A$ then $x \in B$" is equivalent to "$x \in A$ and $x \in B$ if and only if $x \in A$". □

Proposition 2.24. *For any sets A, B and C, we have $C \subseteq A$ if and only if $(A \cap B) \cup C = A \cap (B \cup C)$.*

PROOF: The statement "if $x \in C$ then $x \in A$" is equivalent to the statement "$((x \in A \wedge x \in B) \vee x \in C) \Leftrightarrow (x \in A \wedge (x \in B \vee x \in C))$". □

Definition 2.25 (Disjoint). *Two sets A and B are *disjoint* if and only if $A \cap B = \varnothing$.*

Definition 2.26 (Pairwise Disjoint)**.** Let A be a set. We say the elements of A are *pairwise disjoint* if and only if, for all $x, y \in A$, if $x \cap y \neq \varnothing$ then $x = y$.

Definition 2.27 (Intersection)**.** For any nonempty set $\mathcal{C}$, the *intersection* of $\mathcal{C}$, $\bigcap \mathcal{C}$, is the set that contains exactly those sets that belong to every element of $\mathcal{C}$.

We write $\bigcap_{X \in \mathcal{A}} t[X]$ for $\bigcap \{t[X] \mid X \in \mathcal{A}\}$.

PROOF:
$\langle 1 \rangle 1$. LET: $\mathcal{C}$ be a nonempty set.
$\langle 1 \rangle 2$. There exists a set I whose elements are exactly the sets that belong to every element of $\mathcal{C}$.
   PROOF: Pick $A \in \mathcal{C}$, and take $I = \{x \in A : \forall X \in \mathcal{C}.x \in X\}$.
$\langle 1 \rangle 3$. For any sets I, J, if the elements of I and J are exactly the sets that belong to every element of $\mathcal{C}$ then $I = J$.
   PROOF: Axiom of Extensionality.
$\square$

2.7   Unordered Triples

Definition 2.28 ((Unordered) Triple)**.** Given sets $a_1, \ldots, a_n$, define the *(unordered) n-tuple* $\{a_1, \ldots, a_n\}$ to be

$$\{a_1, \ldots, a_n\} := \{a_1\} \cup \cdots \cup \{a_n\} \ .$$

2.8   Relative Complements

Definition 2.29 (Relative Complement)**.** For any sets A and B, the *difference* or *relative complement* $A - B$ is defined to be

$$A - B := \{x \in A : x \notin B\} \ .$$

Proposition 2.30. *For any sets A and E, we have $A \subseteq E$ if and only if*

$$E - (E - A) = A$$

PROOF:
$\langle 1 \rangle 1$. LET: A and E be sets.
$\langle 1 \rangle 2$. If $A \subseteq E$ then $E - (E - A) = A$
   $\langle 2 \rangle 1$. ASSUME: $A \subseteq E$
   $\langle 2 \rangle 2$. $E - (E - A) \subseteq A$
      PROOF: If $x \in E$ and $x \notin E - A$ then $x \in A$.
   $\langle 2 \rangle 3$. $A \subseteq E - (E - A)$
      PROOF: If $x \in A$ then $x \in E$ and $x \notin E - A$.
$\langle 1 \rangle 3$. If $E - (E - A) = A$ then $A \subseteq E$.
   PROOF: Since $E - (E - A) \subseteq E$.

⬜

Proposition 2.31. *For any set E we have*

$$E - \varnothing = E$$

PROOF: $x \in E$ if and only if $x \in E$ and $x \notin \varnothing$. ⬜

Proposition 2.32. *For any set E we have*

$$E - E = \varnothing \ .$$

PROOF: There is no x such that $x \in E$ and $x \notin E$. ⬜

Proposition 2.33. *For any sets A and E, we have*

$$A \cap (E - A) = \varnothing \ .$$

PROOF: There is no x such that $x \in A$ and $x \in E - A$. ⬜

Proposition 2.34. *Let A and E be sets. Then $A \subseteq E$ if and only if*

$$A \cup (E - A) = E \ .$$

PROOF:
$\langle 1 \rangle 1.$ LET: A and E be sets.
$\langle 1 \rangle 2.$ If $A \subseteq E$ then $A \cup (E - A) = E$.
 $\langle 2 \rangle 1.$ ASSUME: $A \subseteq E$
 $\langle 2 \rangle 2.$ $A \cup (E - A) \subseteq E$
  PROOF: If $x \in A$ or $x \in E - A$ then $x \in E$.
 $\langle 2 \rangle 3.$ $E \subseteq A \cup (E - A)$
  PROOF: If $x \in E$ then either $x \in A$ or $x \notin A$. In the latter case, $x \in E - A$.
$\langle 1 \rangle 3.$ If $A \cup (E - A) = E$ then $A \subseteq E$
 PROOF: Since $A \subseteq A \cup (E - A)$.
⬜

Proposition 2.35. *Let A, B and E be sets. Then:*

 1. If $A \subseteq B$ then $E - B \subseteq E - A$.

 2. If $A \subseteq E$ and $E - B \subseteq E - A$ then $A \subseteq B$.

PROOF:
$\langle 1 \rangle 1.$ LET: A, B and E be sets.
$\langle 1 \rangle 2.$ If $A \subseteq B$ then $E - B \subseteq E - A$.
 PROOF: If $A \subseteq B$, $x \in E$ and $x \notin B$, then we have $x \in E$ and $x \notin A$.
$\langle 1 \rangle 3.$ If $A \subseteq E$ and $E - B \subseteq E - A$ then $A \subseteq B$.
 $\langle 2 \rangle 1.$ ASSUME: $A \subseteq E$
 $\langle 2 \rangle 2.$ ASSUME: $E - B \subseteq E - A$
 $\langle 2 \rangle 3.$ LET: $x \in A$
 $\langle 2 \rangle 4.$ $x \in E$

$\langle 2 \rangle 5.\ x \notin E - A$
$\langle 2 \rangle 6.\ x \notin E - B$
$\langle 2 \rangle 7.\ x \in B$
∎

Example 2.36. We cannot remove the hypothesis $A \subseteq E$ in item 2 above. Let $E = \varnothing$, $A = \{\varnothing\}$ and $B = \varnothing$. Then $E - B = E - A = \varnothing$ but $A \nsubseteq B$.

Proposition 2.37 (De Morgan's Law)**.** *For any sets A, B and E, we have* $E - (A \cup B) = (E - A) \cap (E - B)$.

PROOF: $(x \in E \land \neg(x \in A \lor x \in B)) \Leftrightarrow (x \in E \land x \notin A \land x \in E \land x \notin B)$. ∎

Proposition 2.38 (De Morgan's Law)**.** *For any sets A, B and E, we have* $E - (A \cap B) = (E - A) \cup (E - B)$.

PROOF: $(x \in E \lor \neg(x \in A \land x \in B)) \Leftrightarrow (x \in E \land x \notin A) \lor (x \in E \land x \notin B)$. ∎

Proposition 2.39. *For any sets A, B and E, if $A \subseteq E$ then*

$$A - B = A \cap (E - B)\ .$$

PROOF: If $A \subseteq E$ then we have $(x \in A \land x \notin B) \Leftrightarrow (x \in A \land x \in E \land x \notin B)$. ∎

Proposition 2.40. *For any sets A and B, we have $A \subseteq B$ if and only if $A - B = \varnothing$.*

PROOF: Both are equivalent to the statement that there is no x such that $x \in A$ and $x \notin B$. ∎

Proposition 2.41. *For any sets A and B, we have*

$$A - (A - B) = A \cap B\ .$$

PROOF: $(x \in A \land \neg(x \in A \land x \notin B)) \Leftrightarrow x \in A \land x \in B$. ∎

Proposition 2.42. *For any sets A, B and C, we have*

$$A \cap (B - C) = (A \cap B) - (A \cap C)\ .$$

PROOF: $(x \in A \land x \in B \land x \notin C) \Leftrightarrow (x \in A \land x \in B \land \neg(x \in A \land x \in C))$. ∎

Proposition 2.43. *For any sets A, B, C and E, if $(A \cap B) - C \subseteq E$ then we have*

$$A \cap B \subseteq (A \cap C) \cup (B \cap (E - C))\ .$$

PROOF:
$\langle 1 \rangle 1.$ LET: $x \in A \cap B$
     PROVE: $x \in (A \cap C) \cup (B \cap (E - C))$
$\langle 1 \rangle 2.$ CASE: $x \in C$
  PROOF: Then $x \in A \cap C$.
$\langle 1 \rangle 3.$ CASE: $x \notin C$

8

PROOF: Then $x \in E$ and so $x \in B \cap (E - C)$.
□

Proposition 2.44. *For any sets A, B, C and E, we have*

$$(A \cup C) \cap (B \cup (E - C)) \subseteq A \cup B \ .$$

PROOF: The statement $(x \in A \vee x \in C) \wedge (x \in B \vee (x \in E \wedge x \notin C))$ implies $x \in A \vee x \in B$. □

Proposition 2.45 (De Morgan's Law). *Let E be a set and $\mathcal{C}$ a nonempty set. Then*

$$E - \bigcup \mathcal{C} = \bigcap_{X \in \mathcal{C}} (E - X) \ .$$

PROOF: Easy. □

Proposition 2.46 (De Morgan's Law). *Let E be a set and $\mathcal{C}$ a nonempty set. Then*

$$E - \bigcap \mathcal{C} = \bigcup_{X \in \mathcal{C}} (E - X) \ .$$

PROOF: Easy. □

2.9   Symmetric Difference

Definition 2.47 (Symmetric Difference). For any sets A and B, the *symmetric difference* $A + B$ is defined to be

$$A + B := (A - B) \cup (B - A) \ .$$

Proposition 2.48. *For any sets A and B, we have*

$$A + B = B + A$$

PROOF: From the commutativity of union. □

Proposition 2.49. *For any sets A, B and C, we have*

$$A + (B + C) = (A + B) + C \ .$$

PROOF: Each is the set of all x that belong to either exactly one or all three of A, B and C. □

Proposition 2.50. *For any set A, we have*

$$A + \varnothing = A \ .$$

PROOF:
$$\begin{aligned} A + \varnothing &= (A - \varnothing) \cup (\varnothing - A) \\ &= A \cup \varnothing \\ &= A \end{aligned} \qquad \qquad □$$

Proposition 2.51. *For any set A we have*

$$A + A = \varnothing \ .$$

PROOF:

$$\begin{aligned} A + A &= (A - A) \cup (A - A) \\ &= \varnothing \cup \varnothing \\ &= \varnothing \end{aligned} \qquad \qquad \square$$

2.10 Power Sets

Definition 2.52 (Power Set)**.** For any set A, the *power set* of A, $\mathcal{P}A$, is the set whose elements are exactly the subsets of A.

PROOF: This exists by the Power Set Axiom and Axiom of Comprehension, and is unique by the Axiom of Extensionality. $\square$

Proposition 2.53.

$$\mathcal{P}\varnothing = \{\varnothing\}$$

PROOF: The only subset of $\varnothing$ is $\varnothing$. $\square$

Proposition 2.54. *For any set a, we have*

$$\mathcal{P}\{a\} = \{\varnothing, \{a\}\} \ .$$

PROOF: The only subsets of $\{a\}$ are $\varnothing$ and $\{a\}$. $\square$

Proposition 2.55. *For any sets a and b, we have*

$$\mathcal{P}\{a, b\} = \{\varnothing, \{a\}, \{b\}, \{a, b\}\} \ .$$

PROOF: The only subsets of $\{a, b\}$ are $\varnothing$, $\{a\}$, $\{b\}$ and $\{a, b\}$. $\square$

Proposition 2.56. *For any nonempty set $\mathcal{C}$ we have*

$$\bigcap_{X \in \mathcal{C}} \mathcal{P}X = \mathcal{P}\left(\bigcap \mathcal{C}\right) \ .$$

PROOF:

$$x \in \bigcup_{X \in \mathcal{C}} \mathcal{P}X \Leftrightarrow \forall X \in \mathcal{C}.x \subseteq X$$

$$\begin{aligned} &\Leftrightarrow \forall X \in \mathcal{C}.\forall y \in x.y \in X \\ &\Leftrightarrow \forall y \in x.\forall X \in mathcalC.y \in X \\ &\Leftrightarrow x \subseteq \bigcap \mathcal{C} \end{aligned} \qquad \qquad \square$$

Proposition 2.57. *For any set $\mathcal{C}$ we have*

$$\bigcup_{X \in \mathcal{C}} \mathcal{P}X \subseteq \mathcal{P}\bigcup \mathcal{C} \ .$$

PROOF: If there exists $X \in \mathcal{C}$ such that $x \subseteq X$ then $x \subseteq \bigcup \mathcal{C}$. $\square$

Proposition 2.58. *For any set E, we have*

$$\bigcap \mathcal{P}E = \varnothing \ .$$

PROOF: Since $\varnothing \in \mathcal{P}E$. $\square$

Proposition 2.59. *For any sets E and F, if $E \subseteq F$ then $\mathcal{P}E \subseteq \mathcal{P}F$.*

PROOF: If $E \subseteq F$ and $X \subseteq E$ then $X \subseteq F$. $\square$

Chapter 3

Relations and Functions

3.1 Ordered Pairs

Definition 3.1 (Ordered Pair)**.** For any sets a and b, the *ordered pair* (a, b) is defined by

$$(a, b) := \{\{a\}, \{a, b\}\} \ .$$

Proposition 3.2. *For any sets a, b, x and y, if $(a, b) = (x, y)$ then $a = x$ and $b = y$.*

PROOF:
$\langle 1 \rangle 1$. LET: a, b, x and y be sets.
$\langle 1 \rangle 2$. ASSUME: $(a, b) = (x, y)$
$\langle 1 \rangle 3$. $a = x$
   PROOF: $\{a\} = \bigcap (a, b) = \bigcap (x, y) = \{x\}$.
$\langle 1 \rangle 4$. $\{a, b\} = \{x, y\}$
$\langle 1 \rangle 5$. CASE: $a = b$
   $\langle 2 \rangle 1$. $x = y$
     PROOF: Since $\{x, y\} = \{a, b\}$ is a singleton.
   $\langle 2 \rangle 2$. $b = y$
     PROOF: $b = a = x = y$
$\langle 1 \rangle 6$. CASE: $a \neq b$
   $\langle 2 \rangle 1$. $x \neq y$
     PROOF: Since $\{x, y\} = \{a, b\}$ is not a singleton.
   $\langle 2 \rangle 2$. $b = y$
     PROOF: $\{b\} = \{a, b\} - \{a\} = \{x, y\} - \{x\} = \{y\}$.
□

Definition 3.3 (Cartesian Product)**.** For any sets A and B, the *Cartesian product $A \times B$* is

$$A \times B := \{p \in \mathcal{PP}(A \cup B) : \exists a \in A. \exists b \in B. p = (a, b)\} \ .$$

Proposition 3.4. *For any sets A, B and X, we have*

$$(A - B) \times X = (A \times X) - (B \times X) \ .$$

PROOF: Easy. □

Proposition 3.5. *For any sets A and B, we have $A \times B = \varnothing$ if and only if $A = \varnothing$ or $B = \varnothing$.*

PROOF: Easy. □

Proposition 3.6. *For any sets A, B, X and Y, if $A \subseteq X$ and $B \subseteq Y$ then $A \times B \subseteq X \times Y$. The converse holds assuming $A \neq \varnothing$ and $B \neq \varnothing$.*

PROOF: Easy. □

3.2 Relations

Definition 3.7 (Relation)**.** A *relation* is a set of ordered pairs.
  If R is a relation, we write xRy for $(x, y) \in R$.
  Given sets X and Y, a relation *between X and Y* is a subset of $X \times Y$.
  Given a set X, a relation *on X* is a relation between X and X.

Definition 3.8 (Domain)**.** The *domain* of a relation R is the set

$$\operatorname{dom} R := \left\{ x \in \bigcup\bigcup R : \exists y.(x, y) \in R \right\} \ .$$

Definition 3.9 (Range)**.** The *range* of a relation R is the set

$$\operatorname{ran} R := \left\{ y \in \bigcup\bigcup R : \exists x.(x, y) \in R \right\} \ .$$

Definition 3.10 (Reflexive)**.** Let R be a relation on X. Then R is *reflexive* iff, for all $x \in X$, we have xRx.

Definition 3.11 (Symmetric)**.** Let R be a relation on X. Then R is *symmetric* iff, whenever xRy, then yRx.

Definition 3.12 (Antisymmetric)**.** A relation R is *antisymmetric* iff, whenever xRy and yRx, then $x = y$.

Definition 3.13 (Transitive)**.** Let R be a relation on X. Then R is *transitive* iff, whenever xRy and yRz, then xRz.

Definition 3.14 (Identity Relation)**.** For any set X, the *identity relation I_X* on X is

$$I_X = \{(x, x) : x \in X\} \ .$$

3.3 Composition

Definition 3.15 (Composition)**.** Let R be a relation between X and Y, and S a relation between Y and Z. The *composite* or *relative product* $S \circ R = SR$ is the relation between X and Z defined by

$$x(S \circ R)z \Leftrightarrow \exists y \in Y(xRy \wedge ySz) \ .$$

Proposition 3.16. *Let R be a relation between X and Y, S a relation between Y and Z, and T a relation between Z and W. Then*

$$T(SR) = (TS)R \ .$$

PROOF: Easy. □

Example 3.17. Composition of relations is not commutative in general. Let $X = \{a, b\}$ where $a \neq b$. Let $R = \{(a, a), (b, a)\}$ and $S = \{(a, b), (b, b)\}$. Then $SR = S$ but $RS = R \neq S$.

Proposition 3.18. *A relation R is transitive if and only if $RR \subseteq R$.*

PROOF: Easy. □

3.4 Inverses

Definition 3.19 (Inverse)**.** Let R be a relation between X and Y. The *inverse* or *converse* R^{-1} is the relation between Y and X defined by

$$yR^{-1}x \Leftrightarrow xRy \ .$$

Proposition 3.20. *For any relation R, we have*

$$\operatorname{dom} R^{-1} = \operatorname{ran} R \ .$$

PROOF: Easy. □

Proposition 3.21. *For any relation R, we have*

$$\operatorname{ran} R^{-1} = \operatorname{dom} R \ .$$

PROOF: Easy. □

Proposition 3.22. *Let R be a relation between X and Y, and S a relation between Y and Z. Then*
$$(SR)^{-1} = R^{-1}S^{-1} \ .$$

PROOF: Easy. □

Proposition 3.23. *A relation R is symmetric if and only if $R \subseteq R^{-1}$.*

PROOF: Easy. □

Proposition 3.24. *Let R be a relation between X and Y. Then*

$$I_Y R = R I_X = R \ .$$

PROOF: Easy. $\square$

Proposition 3.25. *A relation R on a set X is reflexive if and only if $I_X \subseteq R$.*

PROOF: Easy. $\square$

Proposition 3.26. *Let R be a relation on a set X. Then R is antisymmetric iff $R \cap R^{-1} \subseteq I_X$.*

PROOF: Easy. $\square$

3.5 Equivalence Relations

Definition 3.27 (Equivalence Relation)**.** Let R be a relation on X. Then R is an *equivalence relation* iff it is reflexive, symmetric and transitive.

Definition 3.28 (Partition)**.** Let X be a set. A *partition* of X is a pairwise disjoint set of nonempty subsets of X whose union is X.

Definition 3.29 (Equivalence Class)**.** Let R be an equivalence relation on X. Let $x \in X$. The *equivalence class* of x with respect to R is

$$x/R := \{y \in X : xRy\} \ .$$

We write X/R for the set of all equivalence classes with respect to R.

Definition 3.30 (Induced)**.** Let P be a partition of X. The relation *induced* by P is X/P where $x(X/P)y$ iff there exists $X \in P$ such that $x \in X$ and $y \in X$.

Theorem 3.31. *Let R be an equivalence relation on X. Then X/R is a partition of X that induces the relation R.*

PROOF: Easy. $\square$

Theorem 3.32. *Let P be a partition of X. Then X/P is an equivalence relation on X, and $P = X/(X/P)$.*

PROOF: Easy. $\square$

3.6 Functions

Definition 3.33 (Function)**.** Let X and Y be sets. A *function, map, mapping, transformation* or *operator f from X to Y*, $f : X \to Y$, is a relation f between X and Y such that, for all $x \in X$, there exists a unique $f(x) \in Y$, called the *value* of f at the *argument x*, such that $(x, f(x)) \in f$.

Definition 3.34 (Onto). Let $f : X \to Y$. We say f maps X *onto* Y iff ran $f = Y$.

Definition 3.35 (Image). Let $f : X \to Y$ and $A \subseteq X$. The *image* of A under f is

$$f(A) := \{f(x) : x \in A\} \ .$$

Definition 3.36 (Inclusion Map). Let Y be a set and $X \subseteq Y$. Then the *inclusion map* $i : X \hookrightarrow Y$ is the function defined by $i(x) = x$ for all $x \in X$.

Proposition 3.37. *For any set X, the identity relation I_X is a function $X \to X$.*

PROOF: Easy. $\square$

Definition 3.38 (Restriction). Let $f : Y \to Z$ and $X \subseteq Y$. The *restriction* of f to X is the function $f{\upharpoonright}X : X \to Z$ defined by

$$(f{\upharpoonright}X)(x) = f(x) \qquad (x \in X) \ .$$

Given sets X, Y and Z with $X \subseteq Y$, if $f : X \to Z$ and $g : Y \to Z$, we say g is an *extension* of f to Y iff $f = g{\upharpoonright}X$.

Definition 3.39 (Projection). Given sets X and Y, the *projection* maps $\pi_1 : X \times Y \to X$ and $\pi_2 : X \times Y \to Y$ are defined by

$$\pi_1(x, y) = x, \qquad \pi_2(x, y) = y \qquad (x \in X, y \in Y) \ .$$

Definition 3.40 (Canonical Map). Let X be a set and R an equivalence relation on X. The *canonical map* $\pi : X \to X/R$ is the map defined by $\pi(x) = x/R$.

Definition 3.41 (One-to-One). A function $f : X \to Y$ is *one-to-one*, or a *one-to-one correspondence*, iff, for all $x, y \in X$, if $f(x) = f(y)$ then $x = y$.

Proposition 3.42. *Let $f : X \to Y$. Then the following are equivalent:*

1. *f is one-to-one.*

2. *For all $A, B \subseteq X$, we have $f(A \cap B) = f(A) \cap f(B)$.*

3. *For all $A \subseteq X$, we have $f(X - A) \subseteq Y - f(A)$.*

PROOF: Easy. $\square$

Proposition 3.43. *Let $f : X \to Y$. Then f maps X onto Y if and only if, for all $A \subseteq X$, we have $Y - f(A) \subseteq f(X - A)$.*

PROOF: Easy. $\square$

3.7 Families

Definition 3.44 (Family)**.** Let I and X be sets. A *family* of elements of X *indexed* by I is a function $a : I \to X$. We write a_i for $a(i)$, and $\{a_i\}_{i \in I}$ for a.

Proposition 3.45 (Generalized Associative Law for Unions)**.** *Let $\{I_j\}_{j \in J}$ be a family of sets. Let $K = \bigcup_{j \in J} I_j$. Let $\{A_k\}_{k \in K}$ be a family of sets indexed by K. Then*

$$\bigcup_{k \in K} A_k = \bigcup_{j \in J} \bigcup_{i \in I_j} A_i \ .$$

PROOF: Easy. $\square$

Proposition 3.46 (Generalized Commutative Law for Unions)**.** *Let $\{I_j\}_{j \in J}$ be a family of sets. Let $f : J \to J$ be a one-to-one correspondence from J onto J. Then*

$$\bigcup_{j \in J} I_j = \bigcup_{j \in J} I_{f(j)} \ .$$

PROOF: Easy. $\square$

Proposition 3.47 (Generalized Associative Law for Intersections)**.** *Let $\{I_j\}_{j \in J}$ be a nonempty family of nonempty sets. Let $K = \bigcup_{j \in J} I_j$. Let $\{A_k\}_{k \in K}$ be a family of sets indexed by K. Then*

$$\bigcap_{k \in K} A_k = \bigcap_{j \in J} \bigcap_{i \in I_j} A_i \ .$$

PROOF: Easy. $\square$

Proposition 3.48 (Generalized Commutative Law for Intersections)**.** *Let $\{I_j\}_{j \in J}$ be a nonempty family of sets. Let $f : J \to J$ be a one-to-one correspondence from J onto J. Then*

$$\bigcap_{j \in J} I_j = \bigcap_{j \in J} I_{f(j)} \ .$$

PROOF: Easy. $\square$

Proposition 3.49. *Let B be a set and $\{A_i\}_{i \in I}$ a family of sets. Then*

$$B \cap \bigcup_{i \in I} A_i = \bigcup_{i \in I} (B \cap A_i)$$

PROOF: Easy. $\square$

Proposition 3.50. *Let B be a set and $\{A_i\}_{i \in I}$ a nonempty family of sets. Then*

$$B \cup \bigcap_{i \in I} A_i = \bigcap_{i \in I} (B \cup A_i)$$

PROOF: Easy. $\square$

Definition 3.51 (Cartesian Product of a Family of Sets)**.** Let $\{A_i\}_{i\in I}$ be a family of sets. The *Cartesian product* $\times_{i\in I} A_i$ is the set of all families $\{a_i\}_{i\in I}$ such that $\forall i \in I.a_i \in A_i$.

We write A^I for $\times_{i\in I} A$.

Definition 3.52 (Projection)**.** Let $\{A_i\}_{i\in I}$ be a family of sets and $i \in I$. The *projection* function $\pi_i : \times_{i\in I} A_i \to A_i$ is defined by $\pi_i(a) = a_i$.

Proposition 3.53. *Let $\{A_i\}_{i\in I}$ and $\{B_j\}_{j\in J}$ be families of sets. Then*

$$\left(\bigcup_{i\in I} A_i\right) \times \left(\bigcup_{j\in J} B_j\right) = \bigcup_{i\in I}\bigcup_{j\in J}(A_i \times B_j) \ .$$

PROOF: Easy. □

Proposition 3.54. *Let $\{A_i\}_{i\in I}$ and $\{B_j\}_{j\in J}$ be nonempty families of sets. Then*

$$\left(\bigcap_{i\in I} A_i\right) \times \left(\bigcap_{j\in J} B_j\right) = \bigcap_{i\in I}\bigcap_{j\in J}(A_i \times B_j) \ .$$

PROOF: Easy. □

Proposition 3.55. *Let $f : X \to Y$. Let $\{A_i\}_{i\in I}$ be a family of subsets of X. Then*

$$f\left(\bigcup_{i\in I} A_i\right) = \bigcup_{i\in I} f(A_i) \ .$$

PROOF: Easy. □

Example 3.56. It is not true in general that, if $f : X \to Y$ and $\{A_i\}_{i\in I}$ is a nonempty family of subsets of X, then $f\left(\bigcap_{i\in I} A_i\right) = \bigcap_{i\in I} f(A_i)$.

Take $X = \{a, b\}$ and $Y = \{c\}$ where $a \neq b$. Take $I = \{i, j\}$ with $i \neq j$. Let $A_i = \{a\}$ and $A_j = \{b\}$. Let f be the unique function $X \to Y$. Then $f\left(\bigcap_{i\in I} A_i\right) = f(\emptyset) = \emptyset$ but $\bigcap_{i\in I} f(A_i) = \{c\}$.

3.8   Inverses and Composites of Functions

Definition 3.57 (Inverse)**.** Given a function $f : X \to Y$, the *inverse* of f is the function $f^{-1} : \mathcal{P}Y \to \mathcal{P}X$ defined by

$$f^{-1}(B) = \{x \in X : f(x) \in B\} \ .$$

We call $f^{-1}(B)$ the *inverse image* of B under f.

Proposition 3.58. *Let $f : X \to Y$. Then f maps X onto Y if and only if the inverse image of any nonempty subset of Y is nonempty.*

PROOF: Easy. □

Proposition 3.59. *Let $f : X \to Y$. Then f is one-to-one if and only if the inverse image of any singleton subset of Y is a singleton.*

PROOF: Easy. □

Proposition 3.60. *Let $f : X \to Y$. Let $B \subseteq Y$. Then*

$$f(f^{-1}(B)) \subseteq B \ .$$

PROOF: Easy. □

Proposition 3.61. *Let $f : X \to Y$. Let $A \subseteq X$. Then*

$$A \subseteq f^{-1}(f(A)) \ .$$

Equality holds if f is one-to-one.

PROOF: Easy. □

Proposition 3.62. *Let $f : X \to Y$. Let $\{B_i\}_{i \in I}$ be a family of subsets of Y. Then*

$$f^{-1}\left(\bigcup_{i \in I} B_i\right) = \bigcup_{i \in I} f^{-1}(B_i) \ .$$

PROOF: Easy. □

Proposition 3.63. *Let $f : X \to Y$. Let $\{B_i\}_{i \in I}$ be a nonempty family of subsets of Y. Then*

$$f^{-1}\left(\bigcap_{i \in I} B_i\right) = \bigcap_{i \in I} f^{-1}(B_i) \ .$$

PROOF: Easy. □

Proposition 3.64. *Let $f : X \to Y$ and $B \subseteq Y$. Then $f^{-1}(Y - B) = X - f^{-1}(B)$.*

PROOF: Easy. □

Proposition 3.65. *Let $f : X \to Y$ be one-to-one. Then the inverse of f as a relation, f^{-1}, is a function $f^{-1} : \operatorname{ran} f \to X$, and for all $y \in \operatorname{ran} f$, we have $f^{-1}(y)$ is the unique x such that $f(x) = y$.*

PROOF: Easy. □

Proposition 3.66. *Let $f : X \to Y$ and $g : Y \to Z$. Then $gf : X \to Z$ and, for all $x \in X$, we have*

$$(g \circ f)(x) = g(f(x)) \ .$$

PROOF: Easy. □

Example 3.67. Example 3.17 shows that function composition is not commutative in general.

Proposition 3.68. *Let* $f : X \to Y$ *and* $g : Y \to Z$. *Then*

$$(gf)^{-1} = f^{-1}g^{-1} : \mathcal{P}Z \to \mathcal{P}X \ .$$

PROOF: Easy. □

Proposition 3.69. *Let* $f : X \to Y$ *and* $g : Y \to X$. *If* $gf = I_X$ *then* f *is one-to-one and* g *maps* Y *onto* X.

PROOF: Easy. □

Chapter 4

Equivalence

Definition 4.1 (Equivalent)**.** Sets E and F are *equivalent*, $E \sim F$, iff there exists a one-to-one correspondence between them.

Proposition 4.2. *For any set X, equivalence is an equivalence relation on $\mathcal{P}X$.*

PROOF: Easy.

Chapter 5

Order

Definition 5.1 (Partial Order)**.** A *partial order* on a set X is a relation on X that is reflexive, antisymmetric and transitive.

A *partially ordered set* or *poset* is a pair $(X, \leqslant)$ such that $\leqslant$ is a partial order on X. We write X for the poset $(X, \leqslant)$.

Given a partial order $\leqslant$, we write $\geqslant$ for the inverse of $\leqslant$.

We write $x < y$ or $y > x$ for $x \leqslant y \wedge x \neq y$. When this holds, we say x is *less than* y, *smaller* than y, or a *predecessor* of y; and y is *greater than* x, *larger* than x, or a *successor* of x.

Proposition 5.2. *For any set X, the relation $\subseteq$ is a partial order on $\mathcal{P}X$.*

PROOF: Easy. □

Proposition 5.3. *In a poset, we never have $x < y$ and $y < x$.*

PROOF: We would then have $x \leqslant y$ and $y \leqslant x$ hence $x = y$ by antisymmetry. But if $x < y$ or $y < x$ then $x \neq y$. □

Proposition 5.4. *The relation $<$ is transitive.*

PROOF:
$\langle 1 \rangle 1$. ASSUME: $x < y$ and $y < z$
$\langle 1 \rangle 2$. $x \leqslant y$ and $y \leqslant z$
$\langle 1 \rangle 3$. $x \leqslant z$
   PROOF: Since $\leqslant$ is transitive.
$\langle 1 \rangle 4$. $x \neq z$
   PROOF: By Proposition 5.3.
□

Proposition 5.5. *Let $<$ be a transitive relation on X such that we never have $x < y$ and $y < x$. Define $\leqslant$ by: $x \leqslant y$ iff $x < y$ or $x = y$. Then $\leqslant$ is a partial order on X.*

PROOF:

$\langle 1 \rangle 1.$ $\leqslant$ is reflexive.

    PROOF: By definition.

$\langle 1 \rangle 2.$ $\leqslant$ is asymmetric.

    PROOF: If $x \leqslant y$ and $y \leqslant x$, we must have $x = y$, because otherwise we would have $x < y$ and $y < x$.

$\langle 1 \rangle 3.$ $\leqslant$ is transitive.

    $\langle 2 \rangle 1.$ LET: $x \leqslant y$ and $y \leqslant z$

    $\langle 2 \rangle 2.$ CASE: $x = y$

      PROOF: We have $y \leqslant z$ so $x \leqslant z$.

    $\langle 2 \rangle 3.$ CASE: $y = z$

      PROOF: We have $x \leqslant y$ so $x \leqslant z$.

    $\langle 2 \rangle 4.$ CASE: $x < y$ and $y < z$

      PROOF: We have $x < z$ by transitivity, so $x \leqslant z$.

□

Definition 5.6 ((Strict) Initial Segment). Let X be a poset and $a \in X$. The *(strict) initial segment* determined by a is

$$s(a) := \{x \in X : x < a\} \ .$$

Definition 5.7 (Weak Initial Segment). Let X be a poset and $a \in X$. The *weak initial segment* determined by a is

$$\bar{s}(a) := \{x \in X : x \leqslant a\} \ .$$

Definition 5.8 (Immediate Successor). Let X be a poset and $x, y \in X$. Then y is the *immediate successor* of x, and x is the *immediate predecessor* of y, iff $x < y$ and there is no z such that $x < z < y$.

Definition 5.9 (Least). Let X be a partial order and $a \in X$. Then a is *least* in X iff $\forall x \in X.a \leqslant x$.

Proposition 5.10. *A poset has at most one least element.*

PROOF: If a and b are least then $a \leqslant b$ and $b \leqslant a$, hence $a = b$. □

Definition 5.11 (Greatest). Let X be a partial order and $a \in X$. Then a is *greatest* in X iff $\forall x \in X.x \leqslant a$.

Proposition 5.12. *A poset has at most one greatest element.*

PROOF: If a and b are greatest then $a \leqslant b$ and $b \leqslant a$, hence $a = b$. □

Definition 5.13 (Total Order). A partial order $\leqslant$ on a set X is a *total order*, *simple order* or *linear order* iff, for all $x, y \in X$, either $x \leqslant y$ or $y \leqslant x$. We then call the poset $(X, \leqslant)$ a *linearly ordered set* or a *chain*.

Proposition 5.14. *Let R be a partial order on X. Then R is total if and only if $X^2 \subseteq R \cup R^{-1}$.*

PROOF: Easy. □

Proposition 5.15. *For any set X, the relation $\subseteq$ is a total order on X iff X is either $\varnothing$ or a singleton.*

PROOF: Easy. □

Chapter 6

Natural Numbers

6.1 Natural Numbers

Definition 6.1 (Successor). The *successor* of a set x, x^+, is defined by

$$x^+ := x \cup \{x\} \ .$$

Definition 6.2. We define

$$0 = \varnothing$$
$$1 = 0^+$$
$$2 = 1^+$$

etc.

Definition 6.3 (Characteristic Function). Let X be a set and $A \subseteq X$. The *characteristic function* of A is the function $\chi_A : X \to 2$ defined by

$$\chi_A(x) = \begin{cases} 1 & \text{if } x \in A \\ 0 & \text{if } x \notin A \end{cases}$$

Theorem 6.4. *Let X be a set. The function $\chi : \mathcal{P}X \to 2^X$ that maps a subset A of X to χ_A is a one-to-one correspondence.*

PROOF: Easy. $\square$

Definition 6.5. The set ω of *natural numbers* is the set such that:

- $0 \in \omega$

- For all $n \in \omega$ we have $n^+ \in \omega$

- For any set X, if $0 \in X$ and $\forall n \in X.n^+ \in X$ then $\omega \subseteq X$

PROOF: To show this exists, pick a set A such that $0 \in A$ and $\forall n \in A.n^+ \in A$ (by the Axiom of Infinity), and let $\omega = \bigcap\{X \in \mathcal{P}A : 0 \in X \wedge \forall n \in X.n^+ \in X\}$. □

Definition 6.6 (Sequence). A *finite sequence* is a family whose index set is a natural number. An *infinite sequence* is a family whose index set is ω.

Given a finite sequence of sets $\{A_i\}_{i \in n^+}$, we write $\bigcup_{i=0}^{n} A_i$ for $\bigcup_{i \in n^+} A_i$. Given an infinite sequence of sets $\{A_i\}_{i \in \omega}$, we write $\bigcup_{i=0}^{\infty} A_i$ for $\bigcup_{i \in \omega} A_i$.

We make similar definitions for $\bigcap$ and $\times$.

Proposition 6.7. *For any natural numbers m and n, if $m \in n$ then $m^+ \in n^+$.*

PROOF:
$\langle 1 \rangle 1$. LET: $P(n)$ be the property $\forall m \in n.m^+ \in n^+$
$\langle 1 \rangle 2$. $P(0)$
  PROOF: Vacuous.
$\langle 1 \rangle 3$. For any natural number n, if $P(n)$ then $P(n^+)$.
  $\langle 2 \rangle 1$. LET: n be a natural number.
  $\langle 2 \rangle 2$. ASSUME: $P(n)$
  $\langle 2 \rangle 3$. LET: $m \in n^+$
  $\langle 2 \rangle 4$. $m \in n$ or $m = n$
  $\langle 2 \rangle 5$. $m^+ \in n^+$ or $m^+ = n^+$
    PROOF: $\langle 2 \rangle 2$
  $\langle 2 \rangle 6$. CASE: $m^+ \in n^{++}$
□

Theorem 6.8 (Principle of Mathematical Induction). *For any subset S of ω, if $0 \in S$ and $\forall n \in S.n^+ \in S$, then $S = \omega$.*

PROOF: From the definition of ω. □

Proposition 6.9.
$$\forall n \in \omega.\forall x \in n.n \nsubseteq x$$

PROOF:
$\langle 1 \rangle 1$. $\forall x \in 0.0 \nsubseteq x$
  PROOF: Vacuous.
$\langle 1 \rangle 2$. For any natural number n, if $\forall x \in n.n \nsubseteq x$ then $\forall x \in n^+.n^+ \nsubseteq x$.
  $\langle 2 \rangle 1$. LET: n be a natural number.
  $\langle 2 \rangle 2$. ASSUME: $\forall x \in n.n \nsubseteq x$
  $\langle 2 \rangle 3$. LET: $x \in n^+$
  $\langle 2 \rangle 4$. ASSUME: for a contradiction $n^+ \subseteq x$
  $\langle 2 \rangle 5$. $x \in n$ or $x = n$
  $\langle 2 \rangle 6$. CASE: $x \in n$
    PROOF: Then we have $n \subseteq n^+ \subseteq x$ contradicting $\langle 2 \rangle 2$.
  $\langle 2 \rangle 7$. CASE: $x = n$
    PROOF: Then we have $n \in n^+ \subseteq x = n$ and $n \subseteq n$ contradicting $\langle 2 \rangle 2$.
□

Corollary 6.9.1. *For any natural number n we have $n \notin n$.*

Corollary 6.9.2. *For any natural number n we have $n \neq n^+$.*

Definition 6.10 (Transitive Set). A set E is a *transitive* set iff, whenever $x \in y \in E$, then $x \in E$.

Proposition 6.11. *Every natural number is a transitive set.*

PROOF:
$\langle 1 \rangle 1$. 0 is a transitive set.
  PROOF: Vacuously, if $x \in y \in 0$ then $x \in 0$.
$\langle 1 \rangle 2$. For any natural number n, if n is a transitive set, then n^+ is a transitive set.
  $\langle 2 \rangle 1$. LET: n be a natural number.
  $\langle 2 \rangle 2$. ASSUME: n is a transitive set.
  $\langle 2 \rangle 3$. LET: $x \in y \in n^+$
  $\langle 2 \rangle 4$. $y \in n$ or $y = n$
  $\langle 2 \rangle 5$. CASE: $y \in n$
    $\langle 3 \rangle 1$. $x \in n$
      PROOF: $\langle 2 \rangle 2$, $\langle 2 \rangle 3$, $\langle 2 \rangle 5$.
    $\langle 3 \rangle 2$. $x \in n^+$
  $\langle 2 \rangle 6$. CASE: $y = n$
    $\langle 3 \rangle 1$. $x \in n$
      PROOF: $\langle 2 \rangle 3$, $\langle 2 \rangle 6$
    $\langle 3 \rangle 2$. $x \in n^+$
□

Proposition 6.12. *For any natural numbers m and n, if $m^+ = n^+$ then $m = n$.*

PROOF:
$\langle 1 \rangle 1$. LET: m and n be natural numbers.
$\langle 1 \rangle 2$. ASSUME: $m^+ = n^+$
$\langle 1 \rangle 3$. $m \in m^+ = n^+$
$\langle 1 \rangle 4$. $m \in n$ or $m = n$
$\langle 1 \rangle 5$. $n \in n^+ = m^+$
$\langle 1 \rangle 6$. $n \in m$ or $n = m$
$\langle 1 \rangle 7$. We cannot have $m \in n$ and $n \in m$
  $\langle 2 \rangle 1$. ASSUME: for a contradiction $m \in n$ and $n \in m$
  $\langle 2 \rangle 2$. $m \in m$
    PROOF: Since m is a transitive set (Proposition 6.11).
  $\langle 2 \rangle 3$. Q.E.D.
    PROOF: This contradicts Proposition 6.9.
$\langle 1 \rangle 8$. $m = n$
□

Theorem 6.13 (Recursion Theorem). *Let X be a set. Let $a \in X$. Let $f : X \to X$. There exists a function $u : \omega \to X$ such that $u(0) = a$ and, for all $n \in \omega$, we have $u(n^+) = f(u(n))$.*

Proof:

⟨1⟩1. LET: $\mathcal{C} = \{A \in \mathcal{P}(\omega \times X) : (0, a) \in A \land \forall n \in \omega.\forall x \in X.(n, x) \in A \Rightarrow$
$\qquad (n^+, f(x)) \in A\}$

⟨1⟩2. $\mathcal{C} \neq \varnothing$
  Proof: $\omega \times X \in \mathcal{C}$

⟨1⟩3. LET: $u = \bigcap \mathcal{C}$

⟨1⟩4. $u \in \mathcal{C}$

⟨1⟩5. u is a function.

  ⟨2⟩1. LET: $P(n)$ be the property: $\forall x, y \in X.(n, x) \in u \land (n, y) \in u \Rightarrow x = y$

  ⟨2⟩2. $P(0)$

    ⟨3⟩1. $\forall x \in X.(0, x) \in u \Rightarrow x = a$
      Proof: If $(0, x) \in u$ and $x \neq a$ then $u - \{(0, x)\} \in \mathcal{C}$ and so $u - \{(0, x)\} \subseteq u$,
      which is impossible.

  ⟨2⟩3. For every natural number n, if $P(n)$ then $P(n^+)$.

    ⟨3⟩1. LET: n be a natural number.

    ⟨3⟩2. ASSUME: $P(n)$

    ⟨3⟩3. LET: $x, y \in X$

    ⟨3⟩4. ASSUME: $(n^+, x), (n^+, y) \in u$

    ⟨3⟩5. PICK $x', y' \in X$ such that $(n, x') \in u$, $(n, y') \in u$ and $f(x') = x$ and
      $f(y') = y$
      Proof: If no such x' exists then $u - \{(n^+, x)\} \in \mathcal{C}$ and so $u - \{(n^+, x)\} \subseteq u$
      which is impossible. Similarly for y'.

    ⟨3⟩6. $x' = y'$
      Proof: ⟨3⟩2

    ⟨3⟩7. $x = y$

□

Proposition 6.14. *For any natural number n, either $n = 0$ or there exists a natural number m such that $n = m^+$.*

Proof: Easy induction on n. □

Proposition 6.15. *ω is a transitive set.*

Proof:

⟨1⟩1. LET: $P(n)$ be the property $\forall x \in n.x \in \omega$

⟨1⟩2. $P(0)$
  Proof: Vacuous.

⟨1⟩3. For any natural number n, if $P(n)$ then $P(n^+)$.

  ⟨2⟩1. LET: n be a natural number.

  ⟨2⟩2. ASSUME: $P(n)$

  ⟨2⟩3. LET: $x \in n^+$

  ⟨2⟩4. $x \in n$ or $x = n$

  ⟨2⟩5. CASE: $x \in n$
    Proof: Then $x \in \omega$ by ⟨2⟩2.

  ⟨2⟩6. CASE: $x = n$
    Proof: Then $x \in \omega$ by ⟨2⟩1.

⬜

Proposition 6.16. *For any natural number n and any nonempty subset $E \subseteq n$, there exists $k \in E$ such that $\forall m \in E . k = m \vee k \in m$.*

PROOF:
$\langle 1 \rangle 1$. LET: $P(n)$ be the property: for any nonempty subset $E \subseteq n$, there exists
$\qquad k \in E$ such that $\forall m \in E . k = m \vee k \in m$
$\langle 1 \rangle 2$. $P(0)$
$\quad$ PROOF: Vacuous as there is no nonempty subset of 0.
$\langle 1 \rangle 3$. For any natural number n, if $P(n)$ then $P(n^+)$.
$\quad \langle 2 \rangle 1$. LET: n be a natural number.
$\quad \langle 2 \rangle 2$. ASSUME: $P(n)$
$\quad \langle 2 \rangle 3$. LET: E be a nonempty subset of n^+
$\quad \langle 2 \rangle 4$. CASE: $E - \{n\} = \varnothing$
$\quad\quad$ PROOF: Then $E = \{n\}$ so take $k = n$.
$\quad \langle 2 \rangle 5$. CASE: $E - \{n\} \neq \varnothing$
$\quad\quad \langle 3 \rangle 1$. PICK $k \in E - \{n\}$ such that $\forall m \in E - \{n\} . k = m \vee k \in m$
$\quad\quad\quad$ PROOF: By $\langle 2 \rangle 2$.
$\quad\quad \langle 3 \rangle 2$. $\forall m \in E . k = m \vee k \in m$
$\quad\quad\quad$ PROOF: Since $k \in n$.
⬜

6.2 Arithmetic

Definition 6.17 (Addition). Define *addition* $+$ on ω by recursion thus:

$$m + 0 = m$$
$$m + n^+ = (m + n)^+$$

Proposition 6.18. *For all $m, n, p \in \omega$ we have*

$$m + (n + p) = (m + n) + p \ .$$

PROOF:
$\langle 1 \rangle 1$. LET: $P(p)$ be the property $\forall m, n \in \omega . m + (n + p) = (m + n) + p$
$\langle 1 \rangle 2$. $P(0)$
$\quad$ PROOF: $m + (n + 0) = m + n = (m + n) + 0$.
$\langle 1 \rangle 3$. $\forall p \in \omega . P(p) \Rightarrow P(p^+)$
$\quad \langle 2 \rangle 1$. LET: $p \in \omega$
$\quad \langle 2 \rangle 2$. ASSUME: $P(p)$
$\quad \langle 2 \rangle 3$. LET: $m, n \in \omega$
$\quad \langle 2 \rangle 4$. $m + (n + p^+) = (m + n) + p^+$

PROOF:

$$m + (n + p^+) = m + (n + p)^+$$
$$= (m + (n + p))^+$$
$$= ((m + n) + p)^+$$
$$= (m + n) + p^+$$

□

Proposition 6.19. *For all* $m, n \in \omega$, *we have*

$$m + n = n + m \ .$$

PROOF:
⟨1⟩1. LET: $P(m)$ be the property $\forall n \in \omega . m + n = n + m$
⟨1⟩2. $P(0)$
  ⟨2⟩1. LET: $Q(n)$ be the property $0 + n = n + 0$
  ⟨2⟩2. $Q(0)$
    PROOF: Trivial.
  ⟨2⟩3. $\forall n \in \omega . Q(n) \Rightarrow Q(n^+)$
    ⟨3⟩1. LET: $n \in \omega$
    ⟨3⟩2. ASSUME: $Q(n)$
    ⟨3⟩3. $0 + n^+ = n^+ + 0$
      PROOF:

$$0 + n^+ = (0 + n)^+$$
$$= (n + 0)^+ \qquad\qquad (\langle 3 \rangle 2)$$
$$= n^+$$
$$= n^+ + 0$$

⟨1⟩3. $\forall m \in \omega . P(m) \Rightarrow P(m^+)$
  ⟨2⟩1. LET: $m \in \omega$
  ⟨2⟩2. ASSUME: $P(m)$
  ⟨2⟩3. LET: $Q(n)$ be the property $m^+ + n = n + m^+$
  ⟨2⟩4. $Q(0)$
    PROOF: ⟨1⟩2
  ⟨2⟩5. $\forall n \in \omega . Q(n) \Rightarrow Q(n^+)$
    ⟨3⟩1. LET: $n \in \omega$
    ⟨3⟩2. ASSUME: $Q(n)$
    ⟨3⟩3. $Q(n^+)$

PROOF:

$$m^+ + n^+ = (m^+ + n)^+$$
$$= (n + m^+)^+ \qquad (\langle 3 \rangle 2)$$
$$= (n + m)^{++}$$
$$= (m + n)^{++} \qquad (\langle 2 \rangle 2)$$
$$= (m + n^+)^+$$
$$= (n^+ + m)^+ \qquad (\langle 2 \rangle 2)$$
$$= n^+ + m^+$$

☐

Definition 6.20 (Multiplication). Define *multiplication* $\cdot$ on ω by

$$m0 = 0$$
$$mn^+ = mn + m$$

Proposition 6.21. *For all $m, n, p \in \omega$, we have*

$$m(n + p) = mn + mp \ .$$

PROOF:
$\langle 1 \rangle 1$. LET: $P(p)$ be the statement $\forall m, n \in \omega . m(n + p) = mn + mp$
$\langle 1 \rangle 2$. $P(0)$
   PROOF:

$$m(n + 0) = mn$$
$$= mn + 0$$
$$= mn + m0$$

$\langle 1 \rangle 3$. $\forall p \in \omega . P(p) \Rightarrow P(p^+)$
  $\langle 2 \rangle 1$. LET: $p \in \omega$
  $\langle 2 \rangle 2$. ASSUME: $P(p)$
  $\langle 2 \rangle 3$. LET: $m, n \in \omega$
  $\langle 2 \rangle 4$. $m(n + p^+) = mn + mp^+$
    PROOF:

$$m(n + p^+) = m(n + p)^+$$
$$= m(n + p) + m$$
$$= (mn + mp) + m \qquad (\langle 2 \rangle 2)$$
$$= mn + (mp + m) \qquad \text{(Proposition 6.18)}$$
$$= mn + mp^+$$

☐

Proposition 6.22. *For all $m, n, p \in \omega$ we have*

$$m(np) = (mn)p \ .$$

PROOF:

$\langle 1 \rangle$1. LET: $P(p)$ be the statement $\forall m, n \in \omega.m(np) = (mn)p$

$\langle 1 \rangle$2. $P(0)$

    PROOF:

$$m(n0) = m0$$
$$= 0$$
$$= (mn)0$$

$\langle 1 \rangle$3. $\forall p \in \omega.P(p) \Rightarrow P(p^+)$

    $\langle 2 \rangle$1. LET: $p \in \omega$

    $\langle 2 \rangle$2. ASSUME: $P(p)$

    $\langle 2 \rangle$3. LET: $m, n \in \omega$

    $\langle 2 \rangle$4. $m(np^+) = (mn)p^+$

      PROOF:

$$m(np^+) = m(np + n)$$
$$= m(np) + mn \qquad \text{(Proposition 6.21)}$$
$$= (mn)p + mn \qquad\qquad (\langle 2 \rangle 2)$$
$$= (mn)p^+$$

$\square$

Proposition 6.23. *For all $m, n \in \omega$, we have*

$$mn = nm \ \ .$$

PROOF:

$\langle 1 \rangle$1. LET: $P(m)$ be the statement $\forall n \in \omega.mn = nm$

$\langle 1 \rangle$2. $P(0)$

    $\langle 2 \rangle$1. LET: $Q(n)$ be the statement $0n = n0$

    $\langle 2 \rangle$2. $Q(0)$

      PROOF: Trivial.

    $\langle 2 \rangle$3. $\forall n \in \omega.Q(n) \Rightarrow Q(n^+)$

      $\langle 3 \rangle$1. LET: $n \in \omega$

      $\langle 3 \rangle$2. ASSUME: $Q(n)$

      $\langle 3 \rangle$3. $Q(n^+)$

        PROOF:

$$0n^+ = 0n + 0$$
$$= 0n$$
$$= n0 \qquad\qquad (\langle 3 \rangle 2)$$
$$= 0$$
$$= n^+0$$

$\langle 1 \rangle$3. $\forall m \in \omega.P(m) \Rightarrow P(m^+)$

    $\langle 2 \rangle$1. LET: $m \in \omega$

    $\langle 2 \rangle$2. ASSUME: $P(m)$

    $\langle 2 \rangle$3. LET: $Q(n)$ be the statement $m^+n = nm^+$

    $\langle 2 \rangle$4. $Q(0)$

      PROOF: $\langle 1 \rangle 2$

$\langle 2 \rangle 5.$ $\forall n \in \omega . Q(n) \Rightarrow Q(n^+)$

$\quad \langle 3 \rangle 1.$ LET: $n \in \omega$

$\quad \langle 3 \rangle 2.$ ASSUME: $Q(n)$

$\quad \langle 3 \rangle 3.$ $Q(n^+)$

$\qquad$ PROOF:

$$
\begin{aligned}
m^+ n^+ &= m^+ n + m^+ \\
&= (m^+ n + m)^+ \\
&= (nm^+ + m)^+ & (\langle 3 \rangle 2) \\
&= (nm + n + m)^+ \\
&= (mn + m + n)^+ & (\langle 2 \rangle 2, \text{Proposition 6.18, Proposition 6.19}) \\
&= (mn^+ + n)^+ \\
&= (n^+ m + n)^+ & (\langle 2 \rangle 2) \\
&= n^+ m + n^+ \\
&= n^+ m^+
\end{aligned}
$$

$\square$

Definition 6.24 (Exponentiation). Define *exponentiation* on ω by recursion:

$$
m^0 = 1
$$
$$
m^{n^+} = m^n m
$$

Proposition 6.25. *For all $m, n, p \in \omega$ we have*

$$
m^{n+p} = m^n m^p \ .
$$

PROOF:

$\langle 1 \rangle 1.$ $m^{n+0} = m^n m^0$

$\quad$ PROOF:

$$
\begin{aligned}
m^{n+0} &= m^n \\
&= m^n 1 \\
&= m^n m^0
\end{aligned}
$$

$\langle 1 \rangle 2.$ If $m^{n+p} = m^n m^p$ then $m^{n+p^+} = m^n m^{p^+}$

$\quad$ PROOF:

$$
\begin{aligned}
m^{n+p^+} &= m^{n+p} m \\
&= m^n m^p m \\
&= m^n m^{p^+}
\end{aligned}
$$

$\square$

Proposition 6.26. *For all $m, n, p \in \omega$ we have*

$$
(m^n)^p = m^{np} \ .
$$

PROOF:

$\langle 1 \rangle 1.$ $(m^n)^0 = m^{n0}$
    PROOF: Both are equal to 1.
$\langle 1 \rangle 2.$ If $(m^n)^p = m^{np}$ then $(m^n)^{p^+} = m^{np^+}$
    PROOF:
$$(m^n)^{p^+} = (m^n)^p m^n$$
$$= m^{np} m^n$$
$$= m^{np+n} \qquad \text{(Proposition 6.25)}$$
$$= m^{np^+}$$

□

6.3   Order on the Natural Numbers

Definition 6.27. Given natural numbers m and n, we write $m < n$ iff $m \in n$.
    We write $m \leqslant n$ iff $m < n \vee m = n$.

Proposition 6.28. *The relation $\leqslant$ is a total order on ω.*

PROOF:
$\langle 1 \rangle 1.$ $\leqslant$ is a partial order on ω.
  $\langle 2 \rangle 1.$ $<$ is transitive.
    PROOF: Proposition 6.11.
  $\langle 2 \rangle 2.$ We never have $m < n$ and $n < m$.
    PROOF: If $m < n$ and $n < m$ then $m < m$ by Proposition 6.11, contradicting
    Corollary 6.9.1.
  $\langle 2 \rangle 3.$ Q.E.D.
$\langle 1 \rangle 2.$ For all $m, n \in \omega$, either $m \leqslant n$ or $n \leqslant m$.
  $\langle 2 \rangle 1.$ LET: $P(n)$ be the statement: $\forall m \in \omega.m \leqslant n \vee n \leqslant m$
  $\langle 2 \rangle 2.$ $P(0)$
    $\langle 3 \rangle 1.$ LET: $Q(m)$ be the statement: $0 \leqslant m$
    $\langle 3 \rangle 2.$ $Q(0)$
      PROOF: Since $0 \leqslant 0$.
    $\langle 3 \rangle 3.$ $\forall m \in \omega.Q(m) \Rightarrow Q(m+1)$
      PROOF: If $0 \leqslant m$ then $0 < m+1$ by transitivity.
  $\langle 2 \rangle 3.$ $\forall n \in \omega.P(n) \Rightarrow P(n+1)$
    $\langle 3 \rangle 1.$ LET: $n \in \omega$
    $\langle 3 \rangle 2.$ ASSUME: $P(n)$
    $\langle 3 \rangle 3.$ $P(n+1)$
      $\langle 4 \rangle 1.$ LET: $Q(m)$ be the property $m \leqslant n+1 \vee n+1 \leqslant m$
      $\langle 4 \rangle 2.$ $Q(0)$
        PROOF: $\langle 2 \rangle 2$
      $\langle 4 \rangle 3.$ $\forall m \in \omega.Q(m) \Rightarrow Q(m+1)$
        $\langle 5 \rangle 1.$ LET: $m \in \omega$
        $\langle 5 \rangle 2.$ ASSUME: $Q(m)$
        $\langle 5 \rangle 3.$ CASE: $m \leqslant n$
          PROOF: Then $m < n+1$

⟨5⟩4. CASE: $n < m$
  PROOF: Then $n + 1 < m + 1$ by Proposition 6.7, so $n + 1 \leqslant m$.
⟨5⟩5. CASE: $n = m$
  PROOF: Then $n + 1 = m + 1$.

□

Proposition 6.29. *For any natural numbers m and n, we have $m \in n$ if and only if $m \subsetneq n$.*

PROOF:
⟨1⟩1. LET: m and n be natural numbers.
⟨1⟩2. If $m \in n$ then $m \subsetneq n$.
  PROOF: Since n is a transitive set, and $m \neq n$ by Corollary 6.9.1.
⟨1⟩3. If $m \subsetneq n$ then $m \in n$.
  ⟨2⟩1. ASSUME: $m \subsetneq n$
  ⟨2⟩2. $n \notin m$
    PROOF: Proposition 6.9.
  ⟨2⟩3. $m \neq n$
  ⟨2⟩4. $m \in n$
    PROOF: Trichotomy.

□

Proposition 6.30. *For natural numbers m, n and k, if $m < n$ then $m + k < n + k$.*

PROOF:
⟨1⟩1. LET: $m, n \in \omega$
⟨1⟩2. ASSUME: $m < n$
⟨1⟩3. $m + 0 < n + 0$
⟨1⟩4. $\forall k \in \omega. m + k < n + k \Rightarrow m + k^+ < n + k^+$
  PROOF: By Proposition 6.7.

□

Proposition 6.31. *For natural numbers m, n and k, if $m < n$ and $k \neq 0$ then $mk < nk$.*

PROOF:
⟨1⟩1. LET: $m, n \in \omega$
⟨1⟩2. ASSUME: $m < n$
⟨1⟩3. $m1 < n1$
⟨1⟩4. For all $k \in \omega$, if $k \neq 0$ and $mk < nk$ then $m(k + 1) < n(k + 1)$
  PROOF:
$$
\begin{aligned}
m(k + 1) &= mk + m \\
&< mk + n &&\text{(Proposition 6.30)} \\
&< nk + n &&\text{(Proposition 6.30)} \\
&= n(k + 1)
\end{aligned}
$$

□

Proposition 6.32. *For any nonempty set of natural numbers E, there exists $k \in E$ such that $\forall m \in E.k \leqslant m$.*

PROOF:
$\langle 1 \rangle 1.$ LET: $E \subseteq \omega$
$\langle 1 \rangle 2.$ ASSUME: there is no $k \in E$ such that $\forall m \in E.k \leqslant m$.
  PROVE: $E = \varnothing$
$\langle 1 \rangle 3.$ $\forall n \in \omega.n \notin E$
 $\langle 2 \rangle 1.$ LET: $P(n)$ be the property: $\forall m < n.m \notin E$
 $\langle 2 \rangle 2.$ $P(0)$
  PROOF: Vacuous.
 $\langle 2 \rangle 3.$ $\forall n \in \omega.P(n) \Rightarrow P(n+1)$
  $\langle 3 \rangle 1.$ LET: $n \in \omega$
  $\langle 3 \rangle 2.$ ASSUME: $\forall m < n.m \notin E$
  $\langle 3 \rangle 3.$ $n \notin E$
   PROOF: From $\langle 1 \rangle 2$.
  $\langle 3 \rangle 4.$ $\forall m < n+1.m \notin E$
$\square$

Proposition 6.33. *Let n be a natural number. Let X be a proper subset of n. Then there exists $m < n$ such that $X \sim m$.*

PROOF:
$\langle 1 \rangle 1.$ LET: $P(n)$ be the property: for every proper subset $X \subsetneq n$, there exists
   $m < n$ such that $X \sim m$.
$\langle 1 \rangle 2.$ $P(0)$
 PROOF: Vacuous.
$\langle 1 \rangle 3.$ $\forall n \in \omega.P(n) \Rightarrow P(n+1)$
 $\langle 2 \rangle 1.$ LET: $n \in \omega$
 $\langle 2 \rangle 2.$ ASSUME: $P(n)$
 $\langle 2 \rangle 3.$ LET: X be a proper subset of $n+1$
 $\langle 2 \rangle 4.$ CASE: $X - \{n\} = n$
  PROOF: Then $X = n$ so $X \sim n < n+1$.
 $\langle 2 \rangle 5.$ CASE: $X - \{n\} \subsetneq n$
  $\langle 3 \rangle 1.$ PICK $m < n$ such that $X - \{n\} \sim m$
  $\langle 3 \rangle 2.$ $X \sim m$ or $X \sim m+1$
   PROOF: If $n \in X$ then $X \sim m+1$. If $n \notin X$ then $X \sim m$.
$\square$

Proposition 6.34. *For every natural number n, we have n is not equivalent to a proper subset of n.*

PROOF:
$\langle 1 \rangle 1.$ LET: $P(n)$ be the property: every one-to-one function $n \to n$ is onto.
$\langle 1 \rangle 2.$ $P(0)$
 PROOF: The only function $0 \to 0$ is $\varnothing$.
$\langle 1 \rangle 3.$ $\forall n \in \omega.P(n) \Rightarrow P(n+1)$
 $\langle 2 \rangle 1.$ LET: $n \in \omega$

⟨2⟩2. ASSUME: $P(n)$

⟨2⟩3. ASSUME: $f : n + 1 \rightarrow n + 1$ is one-to-one.

⟨2⟩4. LET: $g : n \rightarrow n$ be the function
$$g(k) = \begin{cases} f(k) & \text{if } f(k) < n \\ f(n) & \text{if } f(k) = n \end{cases}$$
PROOF: If $k < n$ and $f(k) = n$ then $f(n) < n$ since f is one-to-one.

⟨2⟩5. g is one-to-one.

  ⟨3⟩1. LET: $k, l < n$

  ⟨3⟩2. ASSUME: $g(k) = g(l)$

  ⟨3⟩3. CASE: $f(k) < n$ and $f(l) < n$

    PROOF: Then $f(k) = g(k) = g(l) = f(l)$ so $k = l$ since f is one-to-one.

  ⟨3⟩4. CASE: $f(k) < n$ and $f(l) = n$

    PROOF: Then $f(k) = g(k) = g(l) = f(n)$ contradicting the fact that f is one-to-one.

  ⟨3⟩5. CASE: $f(k) = n$ and $f(l) < n$

    PROOF: Similar.

  ⟨3⟩6. CASE: $f(k) = n$ and $f(l) = n$

    PROOF: Then $k = l$ since f is one-to-one.

⟨2⟩6. g maps n onto n.

  PROOF: ⟨2⟩2

⟨2⟩7. f maps $n + 1$ onto $n + 1$.

  ⟨3⟩1. LET: $l < n + 1$

  ⟨3⟩2. CASE: $l < n$

    ⟨4⟩1. PICK $k < n$ such that $g(k) = l$

    ⟨4⟩2. $f(k) = l$ or $f(n) = l$

  ⟨3⟩3. CASE: $l = n$

    ⟨4⟩1. CASE: $f(n) = n$

      PROOF: Then $l \in \operatorname{ran} f$ as required.

    ⟨4⟩2. CASE: $f(n) < n$

      ⟨5⟩1. PICK $k < n$ such that $g(k) = f(n)$

      ⟨5⟩2. $f(k) = n$

☐

Corollary 6.34.1. *Equivalent natural numbers are equal.*

6.4   Finite Sets

Definition 6.35 (Finite)**.** A set is *finite* iff it is equivalent to a natural number; otherwise, it is *infinite*.

Proposition 6.36. *No finite set is equivalent to one of its proper subsets.*

PROOF: From Proposition 6.34. ☐

Proposition 6.37. ω *is infinite.*

PROOF: Since the function that maps n to $n + 1$ is a one-to-one correspondence between ω and $\omega - \{0\}$. □

Proposition 6.38. *Every subset of a finite set is finite.*

PROOF: Proposition 6.33. □

Definition 6.39 (Number of Elements). For any finite set E, the *number of elements* in E, $\sharp(E)$, is the unique natural number such that $E \sim \sharp(E)$.

Proposition 6.40. *Let E and F be finite sets. If $E \subseteq F$ then $\sharp(E) \leqslant \sharp(F)$.*

PROOF: Proposition 6.33. □

Proposition 6.41. *Let E and F be disjoint finite sets. Then $E \cup F$ is finite and $\sharp(E \cup F) = \sharp(E) \cup \sharp(F)$.*

PROOF:
$\langle 1 \rangle 1$. LET: $P(n)$ be the statement: $n \in \omega$ and for any $m \in \omega$, if $E \sim m$, $F \sim n$
  and $E \cap F = \varnothing$, then $E \cup F \sim m + n$
$\langle 1 \rangle 2$. $P(0)$
  $\langle 2 \rangle 1$. LET: $m \in \omega$
  $\langle 2 \rangle 2$. LET: $E \sim m$ and $F \sim 0$
  $\langle 2 \rangle 3$. $F = \varnothing$
  $\langle 2 \rangle 4$. $E \cup F = E \sim m = m + 0$
$\langle 1 \rangle 3$. $\forall n \in \omega . P(n) \Rightarrow P(n + 1)$
  $\langle 2 \rangle 1$. LET: $n \in \omega$
  $\langle 2 \rangle 2$. ASSUME: $P(n)$
  $\langle 2 \rangle 3$. LET: $m \in \omega$
  $\langle 2 \rangle 4$. LET: $E \sim m$ and $F \sim n + 1$
  $\langle 2 \rangle 5$. ASSUME: $E \cap F = \varnothing$
  $\langle 2 \rangle 6$. PICK $f \in F$
  $\langle 2 \rangle 7$. $F - \{f\} \sim n$
  $\langle 2 \rangle 8$. $E \cap (F - \{f\}) = \varnothing$
  $\langle 2 \rangle 9$. $E \cup (F - \{f\}) \sim m + n$
    PROOF: $\langle 2 \rangle 2$
  $\langle 2 \rangle 10$. $E \cup F \sim m + n + 1$
□

Corollary 6.41.1. *The union of two finite sets is finite.*

PROOF: Since, if E and F are finite, then $E \cup F = (E - F) \cup (E \cap F) \cup (F - E)$ and these are finite and disjoint. □

Proposition 6.42. *If E and F are finite sets then $E \times F$ is finite and $\sharp(E \times F) = \sharp(E)\sharp(F)$.*

PROOF:
$\langle 1 \rangle 1$. LET: $P(n)$ be the statement: $n \in \omega$ and for all $m \in \omega$, if $E \sim m$ and
  $F \sim n$ then $E \times F \sim mn$

$\langle 1\rangle 2.\ P(0)$

    PROOF: If $F \sim 0$ then $F = \varnothing$ so $E \times F = \varnothing \sim 0$.

$\langle 1\rangle 3.\ \forall n \in \omega.P(n) \Rightarrow P(n+1)$

    $\langle 2\rangle 1.$ LET: $n \in \omega$

    $\langle 2\rangle 2.$ ASSUME: $P(n)$

    $\langle 2\rangle 3.$ LET: $m \in \omega$

    $\langle 2\rangle 4.$ ASSUME: $E \sim m$ and $F \sim n+1$

    $\langle 2\rangle 5.$ PICK $f \in F$

    $\langle 2\rangle 6.$ $F - \{f\} \sim n$

    $\langle 2\rangle 7.$ $E \times (F - \{f\}) \sim mn$

    $\langle 2\rangle 8.$ $E \times F = (E \times (F - \{f\})) \cup (E \times \{f\})$

    $\langle 2\rangle 9.$ $E \times \{f\} \sim m$

    $\langle 2\rangle 10.$ $E \times F \sim mn + m$

      PROOF: Proposition 6.41.

☐

Proposition 6.43. *For any finite sets E and F, we have E^F is finite and $\sharp(E^F) = \sharp(E)^{\sharp(F)}$.*

PROOF:

$\langle 1\rangle 1.$ LET: $P(n)$ be the property: $n \in \omega$ and for all $m \in \omega$, if $E \sim m$ and $F \sim n$
        then $E^F \sim m^n$

$\langle 1\rangle 2.\ P(0)$

    PROOF: Since $E^\varnothing = \{\varnothing\} \sim 1$

$\langle 1\rangle 3.\ \forall n \in \omega.P(n) \Rightarrow P(n+1)$

    $\langle 2\rangle 1.$ LET: $n \in \omega$

    $\langle 2\rangle 2.$ ASSUME: $P(n)$

    $\langle 2\rangle 3.$ LET: $m \in \omega$

    $\langle 2\rangle 4.$ LET: $E \sim m$ and $F \sim n+1$

    $\langle 2\rangle 5.$ PICK $f \in F$

    $\langle 2\rangle 6.$ $F - \{f\} \sim n$

    $\langle 2\rangle 7.$ LET: $\phi : E^F \to E^{F-\{f\}} \times E$ be the function $\phi(g) = (g{\upharpoonright}(F - \{f\}), g(f))$

    $\langle 2\rangle 8.$ ϕ is a one-to-one correspondence

    $\langle 2\rangle 9.$ $\sharp(E^F) = m^{n+1}$

      PROOF:

$$\begin{aligned}
\sharp(E^F) &= \sharp(E^{F-\{f\}} \times E) \\
&= \sharp(E^{F-\{f\}})\sharp(E) && \text{(Proposition 6.42)} \\
&= m^n m && (\langle 2\rangle 2,\ \langle 2\rangle 4) \\
&= m^{n+1}
\end{aligned}$$

☐

Corollary 6.43.1. *If E is finite then $\mathcal{P}E$ is finite and $\sharp(\mathcal{P}E) = 2^{\sharp(E)}$.*

Proposition 6.44. *The union of a finite set of finite sets is finite.*

PROOF:

⟨1⟩1. LET: $P(n)$ be the property: for any set E, if $E \sim n$ and every element of E is finite, then $\bigcup E$ is finite.

⟨1⟩2. $P(0)$
  PROOF: Since $\bigcup \varnothing = \varnothing$ is finite.

⟨1⟩3. $\forall n \in \omega. P(n) \Rightarrow P(n+1)$
  ⟨2⟩1. LET: n be a natural number.
  ⟨2⟩2. ASSUME: $P(n)$
  ⟨2⟩3. LET: $E \sim n + 1$
  ⟨2⟩4. PICK $X \in E$
  ⟨2⟩5. $E - \{X\} \sim n$
  ⟨2⟩6. $\bigcup(E - \{X\})$ is finite.
    PROOF: ⟨2⟩2
  ⟨2⟩7. $\bigcup E = \bigcup(E - \{X\}) \cup X$
  ⟨2⟩8. $\bigcup E$ is finite.
    PROOF: Corollary 6.41.1.
▯

Proposition 6.45. *Every nonempty finite set of natural numbers has a greatest element.*

PROOF:

⟨1⟩1. LET: $P(n)$ be the property: for every $E \subseteq \mathbb{N}$, if $E \sim n$ then E has a greatest element.

⟨1⟩2. $P(1)$
  PROOF: Since k is the greatest element of $\{k\}$.

⟨1⟩3. $\forall n \geqslant 1. P(n) \Rightarrow P(n+1)$
  ⟨2⟩1. LET: $n \geqslant 1$
  ⟨2⟩2. ASSUME: $P(n)$
  ⟨2⟩3. ASSUME: $E \subseteq \omega$ and $E \sim n + 1$
  ⟨2⟩4. PICK $k \in E$
  ⟨2⟩5. LET: l be the greatest element of $E - \{k\}$
  ⟨2⟩6. Either k or l is greatest in E.
▯