

Mathematics

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Chapter 1

Sets and Classes

1.1 Classes

Our language is the language of first-order logic with equality over one primitive binary predicate \in . We call all the objects we reason about *sets*. When $a \in b$, we say a is a *member* or *element* of b , or b *contains* a . We write $b \ni a$ for $a \in b$, and $a \notin b$ for $\neg(a \in b)$. We write $\forall x \in a. \phi$ as an abbreviation for $\forall x(x \in a \rightarrow \phi)$, and $\exists x \in a. \phi$ as an abbreviation for $\exists x(x \in a \wedge \phi)$.

We shall speak informally of *classes* as an abbreviation for talking about predicates. A *class* is determined by a unary predicate $\phi[x]$ (possibly with parameters). We write $\{x \mid \phi[x]\}$ or $\{x : \phi[x]\}$ for the class determined by $\phi[x]$. We write ' a is an element of $\{x \mid \phi[x]\}$ ' or ' $a \in \{x \mid \phi[x]\}$ ' for $\phi[a]$.

We write $\{t[x_1, \dots, x_n] \mid P[x_1, \dots, x_n]\}$ for

$$\{y \mid \exists x_1, \dots, x_n (y = t[x_1, \dots, x_n] \wedge P[x_1, \dots, x_n])\} .$$

We say two classes \mathbf{A} and \mathbf{B} are *equal*, and write $\mathbf{A} = \mathbf{B}$, iff $\forall x(x \in \mathbf{A} \leftrightarrow x \in \mathbf{B})$.

Proposition Schema 1.1.1. *For any class \mathbf{A} , the following is a theorem.*

$$\mathbf{A} = \mathbf{A}$$

PROOF: We have $\forall x(x \in \mathbf{A} \leftrightarrow x \in \mathbf{A})$. \square

Proposition Schema 1.1.2. *For any classes \mathbf{A} and \mathbf{B} , the following is a theorem.*

If $\mathbf{A} = \mathbf{B}$ then $\mathbf{B} = \mathbf{A}$.

PROOF: If $\forall x(x \in \mathbf{A} \leftrightarrow x \in \mathbf{B})$ then $\forall x(x \in \mathbf{B} \leftrightarrow x \in \mathbf{A})$. \square

Proposition Schema 1.1.3. *For any classes \mathbf{A} , \mathbf{B} and \mathbf{C} , the following is a theorem.*

If $\mathbf{A} = \mathbf{B}$ and $\mathbf{B} = \mathbf{C}$ then $\mathbf{A} = \mathbf{C}$.

PROOF: If $\forall x(x \in \mathbf{A} \leftrightarrow x \in \mathbf{B})$ and $\forall x(x \in \mathbf{B} \leftrightarrow x \in \mathbf{C})$ then $\forall x(x \in \mathbf{A} \leftrightarrow x \in \mathbf{C})$. \square

1.1.1 Subclasses

Definition 1.1.4 (Subclass). We say a class \mathbf{A} is a *subclass* of \mathbf{B} , or \mathbf{B} is a *superclass* of \mathbf{A} , or \mathbf{B} *includes* \mathbf{A} , and write $\mathbf{A} \subseteq \mathbf{B}$ or $\mathbf{B} \supseteq \mathbf{A}$, iff every element of \mathbf{A} is an element of \mathbf{B} . Otherwise we write $\mathbf{A} \not\subseteq \mathbf{B}$ or $\mathbf{B} \not\supseteq \mathbf{A}$.

We say \mathbf{A} is a *proper* subclass of \mathbf{B} , \mathbf{B} is a *proper* superclass of \mathbf{A} , or \mathbf{B} *properly* includes \mathbf{A} , and write $\mathbf{A} \subsetneq \mathbf{B}$ or $\mathbf{B} \supsetneq \mathbf{A}$, iff $\mathbf{A} \subseteq \mathbf{B}$ and $\mathbf{A} \neq \mathbf{B}$.

Proposition Schema 1.1.5. *For any class \mathbf{A} , the following is a theorem.*

$$\mathbf{A} \subseteq \mathbf{A}$$

PROOF: Every element of \mathbf{A} is an element of \mathbf{A} . \square

Proposition Schema 1.1.6. *For any classes \mathbf{A} and \mathbf{B} , the following is a theorem.*

If $\mathbf{A} \subseteq \mathbf{B}$ and $\mathbf{B} \subseteq \mathbf{A}$ then $\mathbf{A} = \mathbf{B}$.

PROOF: If every element of \mathbf{A} is an element of \mathbf{B} , and every element of \mathbf{B} is an element of \mathbf{A} , then \mathbf{A} and \mathbf{B} have exactly the same elements. \square

Proposition Schema 1.1.7. *For any classes \mathbf{A} , \mathbf{B} and \mathbf{C} , the following is a theorem.*

If $\mathbf{A} \subseteq \mathbf{B}$ and $\mathbf{B} \subseteq \mathbf{C}$ then $\mathbf{A} \subseteq \mathbf{C}$.

PROOF: If every element of \mathbf{A} is an element of \mathbf{B} and every element of \mathbf{B} is an element of \mathbf{C} then every element of \mathbf{A} is an element of \mathbf{C} . \square

1.1.2 Constructions of Classes

Definition 1.1.8 (Empty Class). The *empty class* \emptyset is $\{x \mid \perp\}$. Every other class is *nonempty*.

Definition 1.1.9 (Universal Class). The *universal class* \mathbf{V} is $\{x \mid \top\}$.

Definition 1.1.10 (Enumeration). Given objects a_1, \dots, a_n , we define the class $\{a_1, \dots, a_n\}$ to be the class $\{x \mid x = a_1 \vee \dots \vee x = a_n\}$.

Definition 1.1.11 (Intersection). For any classes \mathbf{A} and \mathbf{B} , the *intersection* $\mathbf{A} \cap \mathbf{B}$ is $\{x \mid x \in \mathbf{A} \wedge x \in \mathbf{B}\}$.

Definition 1.1.12 (Union). For any classes \mathbf{A} and \mathbf{B} , the *union* $\mathbf{A} \cup \mathbf{B}$ is $\{x \mid x \in \mathbf{A} \vee x \in \mathbf{B}\}$.

Definition 1.1.13 (Relative Complement). Let \mathbf{A} and \mathbf{B} be classes. The *relative complement* of \mathbf{B} in \mathbf{A} is the class $\mathbf{A} - \mathbf{B} := \{x \in \mathbf{A} \mid x \notin \mathbf{B}\}$.

Definition 1.1.14 (Symmetric Difference). For any classes \mathbf{A} and \mathbf{B} , the *symmetric difference* is the class $\mathbf{A} + \mathbf{B} := (\mathbf{A} - \mathbf{B}) \cup (\mathbf{B} - \mathbf{A})$.

Definition 1.1.15 (Pairwise disjoint). Let \mathbf{A} be a class. We say the elements of \mathbf{A} are *pairwise disjoint* iff, for all $x, y \in \mathbf{A}$, if $x \cap y \neq \emptyset$ then $x = y$.

1.2 Sets and the Axiom of Extensionality

Definition 1.2.1 (Axiom of Extensionality). The *Axiom of Extensionality* is the statement: if two sets have exactly the same members, then they are equal.

$$\forall x, y (\forall z (z \in x \Leftrightarrow z \in y) \Rightarrow x = y) .$$

When working in a theory with the Axiom of Extensionality, we may identify a set a with the class $\{x \mid x \in a\}$. Our use of the symbols \in and $=$ is consistent. We say a class \mathbf{A} is a set iff there exists a set a such that $a = \mathbf{A}$; that is, $\{x \mid \phi[x]\}$ is a set iff $\exists a \forall x (x \in a \Leftrightarrow \phi[x])$. Otherwise, \mathbf{A} is a *proper class*.

Definition 1.2.2 (Subset). If A is a set and $A \subseteq \mathbf{B}$, we say A is a *subset* of \mathbf{B} .

Definition 1.2.3 (Union). The *union* of a class \mathbf{A} is $\{x \mid \exists X \in \mathbf{A}. x \in X\}$. We write $\bigcup_{P(x)} t(x)$ for $\bigcup \{t(x) \mid P(x)\}$.

Definition 1.2.4 (Intersection). The *intersection* of a class \mathbf{A} is $\{x \mid \forall X \in \mathbf{A}. x \in X\}$. We write $\bigcap_{P(x)} t(x)$ for $\bigcap \{t(x) \mid P(x)\}$.

Definition 1.2.5 (Power Class). For any class \mathbf{A} , the *power class* $\mathcal{P}\mathbf{A}$ is $\{X \mid X \subseteq \mathbf{A}\}$.

1.3 The Other Axioms

Definition 1.3.1 (Pairing Axiom). The *Pairing Axiom* is the statement: for any sets a and b , the class $\{a, b\}$ is a set.

$$\forall a \forall b \exists c \forall x (x \in c \Leftrightarrow x = a \vee x = b)$$

Definition 1.3.2 (Union Axiom). The *Union Axiom* is the statement: for any set A , the class $\bigcup A$ is a set.

$$\forall A \exists B \forall x (x \in B \Leftrightarrow \exists y (y \in A \wedge x \in y))$$

Definition 1.3.3 (Comprehension Axiom Scheme). The *Comprehension Axiom Scheme* is the set of sentences of the form, for any class \mathbf{A} : If \mathbf{A} is a subclass of a set then \mathbf{A} is a set.

That is, for any property $P[x, y_1, \dots, y_n]$:

For any sets a_1, \dots, a_n and B , the class $\{x \in B \mid P[x, a_1, \dots, a_n]\}$ is a set.

$$\forall a_1, \dots, a_n, B. \exists C. \forall x (x \in C \Leftrightarrow x \in B \wedge P[x, a_1, \dots, a_n])$$

Definition 1.3.4 (Replacement Axiom Scheme). The *Replacement Axiom Scheme* is the set of sentences of the form, for some property $P[x, y, z_1, \dots, z_n]$:

For any sets a_1, \dots, a_n, B , assume for all $x \in B$ there exists at most one y such that $P[x, y, a_1, \dots, a_n]$. Then $\{y \mid \exists x \in B. P[x, y, a_1, \dots, a_n]\}$ is a set.

$$\forall a_1, \dots, a_n, B (\forall x \in B. \forall y, y' (P[x, y, a_1, \dots, a_n] \wedge P[x, y', a_1, \dots, a_n] \Rightarrow y = y') \Rightarrow \\ \exists C \forall y (y \in C \Leftrightarrow \exists x \in B. P[x, y, a_1, \dots, a_n]))$$

Definition 1.3.5 (Power Set Axiom). The *Power Set Axiom* is the statement: the power class of a set is a set.

$$\forall A \exists B \forall x (x \in B \Leftrightarrow \forall y (y \in x \Rightarrow y \in A))$$

Definition 1.3.6 (Axiom of Infinity). The *Axiom of Infinity* is the statement: there exists a set I such that $\emptyset \in I$ and $\forall x \in I. x \cup \{x\} \in I$.

$$\exists I (\emptyset \in I. \forall x. x \notin I \wedge \forall x \in I. \exists y \in I. \forall z (z \in y \Leftrightarrow z \in x \vee z = x))$$

Definition 1.3.7 (Axiom of Choice). The *Axiom of Choice* is the statement: For any set A of pairwise disjoint, nonempty sets, there exists a set C such that, for all $x \in A$, we have $x \cap C$ has exactly one element.

$$\begin{aligned} & \forall A (\forall x \in A. \exists y y \in x \wedge \\ & \forall x, y \in A. \forall z (z \in x \wedge z \in y \Rightarrow x = y) \Rightarrow \\ & \exists C. \forall x \in A. \exists y \forall z (z \in x \wedge z \in C \Leftrightarrow z = y)) \end{aligned}$$

Definition 1.3.8 (Axiom of Regularity). The *Axiom of Regularity* is the statement: for any A , if A has a member, then there exists $m \in A$ such that $m \cap A = \emptyset$.

$$\forall A (\exists x. x \in A \Rightarrow \exists m \in A. \neg \exists x (x \in m \wedge x \in A))$$

Definition 1.3.9 (Zermelo Set Theory). *Zermelo set theory* is the theory whose axioms are:

- Extensionality
- Pairing
- Union
- Comprehension
- Power Set
- Infinity
- Choice
- Regularity

We label theorems with Z when they are provable in Zermelo set theory.

Definition 1.3.10 (Zermelo-Fraenkel Set Theory). *Zermelo-Fraenkel set theory* is the theory whose axioms are:

- Extensionality
- Union

- Replacement
- Power Set
- Infinity
- Choice
- Regularity

We label theorems with ZFC when they are provable in Zermelo-Fraenkel set theory.

We label a theorem with FOL if it can be proved in first-order logic, i.e. from no axioms.

1.4 ZFC Extends Z

Proposition 1.4.1 (Z,ZFC). *The empty class \emptyset is a set.*

PROOF: Immediate from the Axiom of Infinity. \square

Proposition 1.4.2 (ZFC). *The Axiom of Pairing is a theorem of ZFC.*

PROOF:

- $\langle 1 \rangle 1$. LET: a, b be sets.
 $\langle 1 \rangle 2$. LET: $P(x, y)$ be the predicate $(x = \emptyset \wedge y = a) \vee (x = \mathcal{P}\emptyset \wedge y = b)$.
 $\langle 1 \rangle 3$. For all $x \in \mathcal{P}\mathcal{P}\emptyset$, there exists at most one y such that $P(x, y)$.
 $\langle 2 \rangle 1$. LET: $x \in \mathcal{P}\mathcal{P}\emptyset$
 $\langle 2 \rangle 2$. LET: y and y' be sets.
 $\langle 2 \rangle 3$. ASSUME: $P(x, y)$ and $P(x, y')$
 $\langle 2 \rangle 4$. $(x = \emptyset \wedge y = a) \vee (x = \mathcal{P}\emptyset \wedge y = b)$
PROOF: From $\langle 2 \rangle 3$.
 $\langle 2 \rangle 5$. $(x = \emptyset \wedge y' = a) \vee (x = \mathcal{P}\emptyset \wedge y' = b)$
PROOF: From $\langle 2 \rangle 3$.
 $\langle 2 \rangle 6$. $\emptyset \neq \mathcal{P}\emptyset$
PROOF: Since $\emptyset \in \mathcal{P}\emptyset$ and $\emptyset \notin \emptyset$.
 $\langle 2 \rangle 7$. $y = y'$
 $\langle 1 \rangle 4$. LET: A be the set $\{y \mid \exists x \in \mathcal{P}\mathcal{P}\emptyset. P(x, y)\}$.
 $\langle 1 \rangle 5$. $A = \{a, b\}$
 \square

Proposition Schema 1.4.3 (ZFC). *Every instance of the Comprehension Axiom Scheme is a theorem of ZFC.*

PROOF:

- $\langle 1 \rangle 1$. LET: $P(x)$ be a predicate.
 $\langle 1 \rangle 2$. LET: A be a set.
 $\langle 1 \rangle 3$. LET: $Q(x, y)$ be the predicate $P(x) \wedge y = x$.

- ⟨1⟩4. For all $x \in A$, there exists at most one y such that $Q(x, y)$.
 ⟨2⟩1. LET: $x \in A$
 ⟨2⟩2. LET: y and y' be sets.
 ⟨2⟩3. ASSUME: $Q(x, y)$ and $Q(x, y')$
 ⟨2⟩4. $x \in A \wedge P(x) \wedge y = x \wedge y' = x$
 PROOF: From ⟨2⟩3.
 ⟨2⟩5. $y = y'$
 PROOF: From ⟨2⟩4.
 ⟨1⟩5. LET: B be the set $\{y \mid \exists x \in A. Q(x, y)\}$
 PROOF: This is a set by an Axiom of Replacement and ⟨1⟩4.
 ⟨1⟩6. $B = \{y \in A \mid P(y)\}$
 PROOF:

$$\begin{aligned}
 y \in B &\Leftrightarrow \exists x \in A. Q(x, y) && (\langle 1 \rangle 5) \\
 &\Leftrightarrow \exists x \in A (P(x) \wedge y = x) && (\langle 1 \rangle 3) \\
 &\Leftrightarrow P(y)
 \end{aligned}$$

□

Corollary Schema 1.4.3.1 (ZFC). *Every axiom of Z is a theorem of ZFC.*

It follows that every theorem of Z is a theorem of ZFC.

1.5 Consequences of the Axioms

Proposition 1.5.1 (Z). *The union of two sets is a set.*

PROOF: Because $A \cup B = \bigcup \{A, B\}$. □

Proposition Schema 1.5.2 (Z). *For any number n , the following is a theorem:*

For any sets a_1, \dots, a_n , the class $\{a_1, \dots, a_n\} = \{x \mid x = a_1 \vee \dots \vee x = a_n\}$ is a set.

PROOF: The case $n = 1$ follows from Pairing since $\{a\} = \{a, a\}$.

If we have proved the theorem for n we have $\{a_1, \dots, a_n, a_{n+1}\} = \{a_1, \dots, a_n\} \cup \{a_{n+1}\}$. □

Proposition 1.5.3 (Z). *No set is a member of itself.*

PROOF:

- ⟨1⟩1. LET: x be any set.
 ⟨1⟩2. PICK $m \in \{x\}$ such that $m \cap \{x\} = \emptyset$.
 PROOF: Axiom of Regularity.
 ⟨1⟩3. $m = x$
 ⟨1⟩4. $x \cap \{x\} = \emptyset$
 ⟨1⟩5. $x \notin x$

□

Corollary 1.5.3.1 (Z). *The universal class \mathbf{V} is a proper class.*

PROOF: If \mathbf{V} is a set then $\mathbf{V} \in \mathbf{V}$, contradicting the Proposition. \square

Proposition 1.5.4 (Z). *There are no sets a and b such that $a \in b$ and $b \in a$.*

PROOF:

$\langle 1 \rangle 1$. LET: a and b be any sets.

$\langle 1 \rangle 2$. PICK $m \in \{a, b\}$ such that $m \cap \{a, b\} = \emptyset$

$\langle 1 \rangle 3$. CASE: $m = a$

PROOF: Then $b \notin a$.

$\langle 1 \rangle 4$. CASE: $m = b$

PROOF: Then $a \notin b$.

\square

Proposition 1.5.5 (Z). *The intersection of a set and a class is a set.*

PROOF: Immediate from Comprehension. \square

Proposition 1.5.6 (Z). *The relative complement of a class in a set is a set.*

[Z]

PROOF: Immediate from Comprehension. \square

Corollary 1.5.6.1 (Z). *The symmetric difference of two sets is a set.*

Proposition 1.5.7 (Z). *The intersection of a nonempty class is a set.*

PROOF:

$\langle 1 \rangle 1$. LET: \mathbf{A} be a nonempty class.

$\langle 1 \rangle 2$. PICK $B \in \mathbf{A}$

$\langle 1 \rangle 3$. $\bigcap \mathbf{A} \subseteq B$

$\langle 1 \rangle 4$. $\bigcap \mathbf{A}$ is a set.

PROOF: By Comprehension.

\square

Proposition Schema 1.5.8 (FOL). *For any classes \mathbf{A} and \mathbf{B} , the following is a theorem:*

If $\mathbf{A} \subseteq \mathbf{B}$ then $\mathcal{P}\mathbf{A} \subseteq \mathcal{P}\mathbf{B}$.

PROOF: Every subset of \mathbf{A} is a subset of \mathbf{B} . \square

Proposition Schema 1.5.9 (FOL). *For any classes \mathbf{A} and \mathbf{B} , the following is a theorem:*

If $\mathbf{A} \subseteq \mathbf{B}$ then $\bigcup \mathbf{A} \subseteq \bigcup \mathbf{B}$.

PROOF: If $x \in X \in \mathbf{A}$ then $x \in X \in \mathbf{B}$. \square

Proposition Schema 1.5.10 (Z). *For any class \mathbf{A} , the following is a theorem:*

$$\mathbf{A} = \bigcup \mathcal{P}\mathbf{A}$$

PROOF:

$\langle 1 \rangle 1. \mathbf{A} \subseteq \bigcup \mathcal{P}\mathbf{A}$

PROOF: For all $x \in \mathbf{A}$ we have $x \in \{x\} \in \mathcal{P}\mathbf{A}$.

$\langle 1 \rangle 2. \bigcup \mathcal{P}\mathbf{A} \subseteq \mathbf{A}$

$\langle 2 \rangle 1. \text{ LET: } x \in \bigcup \mathcal{P}\mathbf{A}$

$\langle 2 \rangle 2. \text{ PICK } X \in \mathcal{P}\mathbf{A} \text{ such that } x \in X$

$\langle 2 \rangle 3. X \subseteq \mathbf{A}$

$\langle 2 \rangle 4. x \in \mathbf{A}$

□

1.6 Transitive Classes

Definition 1.6.1 (Transitive Class). A class \mathbf{A} is a *transitive class* iff whenever $x \in y \in \mathbf{A}$ then $x \in \mathbf{A}$.

Proposition Schema 1.6.2 (FOL). *For any class \mathbf{A} , the following is a theorem:*

The following are equivalent.

1. \mathbf{A} is a transitive class.

2. $\bigcup \mathbf{A} \subseteq \mathbf{A}$

3. Every element of \mathbf{A} is a subset of \mathbf{A} .

4. $\mathbf{A} \subseteq \mathcal{P}\mathbf{A}$

PROOF: Immediate from definitions. □

Proposition Schema 1.6.3 (FOL). *For any class \mathbf{A} , the following is a theorem:*

If \mathbf{A} is a transitive class then $\bigcup \mathbf{A}$ is a transitive class.

PROOF:

$\langle 1 \rangle 1. \text{ ASSUME: } \mathbf{A} \text{ is a transitive class.}$

$\langle 1 \rangle 2. \text{ LET: } x \in y \in \bigcup \mathbf{A}$

$\langle 1 \rangle 3. y \in \mathbf{A}$

PROOF: Since $\bigcup \mathbf{A} \subseteq \mathbf{A}$ by Proposition 1.6.2.

$\langle 1 \rangle 4. x \in \bigcup \mathbf{A}$

□

Proposition Schema 1.6.4 (Z). *For any class \mathbf{A} , the following is a theorem:*

We have \mathbf{A} is a transitive class if and only if $\mathcal{P}\mathbf{A}$ is a transitive class.

PROOF:

$\langle 1 \rangle 1. \text{ If } \mathbf{A} \text{ is a transitive class then } \mathcal{P}\mathbf{A} \text{ is a transitive class.}$

$\langle 2 \rangle 1. \text{ ASSUME: } \mathbf{A} \text{ is a transitive class.}$

$\langle 2 \rangle 2. \mathbf{A} \subseteq \mathcal{P}\mathbf{A}$

PROOF: Proposition 1.6.2.

$\langle 2 \rangle 3. \mathcal{P}\mathbf{A} \subseteq \mathcal{P}\mathcal{P}\mathbf{A}$

PROOF: Proposition 1.5.8.

$\langle 2 \rangle 4$. $\mathcal{P}\mathbf{A}$ is a transitive class.

PROOF: Proposition 1.6.2.

$\langle 1 \rangle 2$. If $\mathcal{P}\mathbf{A}$ is a transitive class then \mathbf{A} is a transitive class.

$\langle 2 \rangle 1$. ASSUME: $\mathcal{P}\mathbf{A}$ is a transitive class.

$\langle 2 \rangle 2$. $\bigcup \mathcal{P}\mathbf{A} \subseteq \mathcal{P}\mathbf{A}$

PROOF: Proposition 1.6.2.

$\langle 2 \rangle 3$. $\mathbf{A} \subseteq \mathcal{P}\mathbf{A}$

PROOF: Proposition 1.5.10.

$\langle 2 \rangle 4$. \mathbf{A} is a transitive class.

PROOF: Proposition 1.6.2.

□

Proposition Schema 1.6.5 (FOL). *For any class \mathbf{A} , the following is a theorem:*

If every member of \mathbf{A} is a transitive set then $\bigcup \mathbf{A}$ is a transitive class.

PROOF:

$\langle 1 \rangle 1$. ASSUME: Every member of \mathbf{A} is a transitive set.

$\langle 1 \rangle 2$. LET: $x \in y \in \bigcup \mathbf{A}$

$\langle 1 \rangle 3$. PICK $A \in \mathbf{A}$ such that $y \in A$.

$\langle 1 \rangle 4$. $x \in A$

PROOF: Since A is a transitive set.

$\langle 1 \rangle 5$. $x \in \bigcup \mathbf{A}$

□

Proposition Schema 1.6.6 (FOL). *For any class \mathbf{A} , the following is a theorem:*

If every member of \mathbf{A} is a transitive set then $\bigcap \mathbf{A}$ is a transitive class.

PROOF:

$\langle 1 \rangle 1$. ASSUME: Every member of \mathbf{A} is a transitive set.

$\langle 1 \rangle 2$. LET: $x \in y \in \bigcap \mathbf{A}$

PROVE: $x \in \bigcap \mathbf{A}$

$\langle 1 \rangle 3$. LET: $A \in \mathbf{A}$

$\langle 1 \rangle 4$. $y \in A$

$\langle 1 \rangle 5$. $x \in A$

PROOF: Since A is a transitive set.

□

Chapter 2

Relations

2.1 Ordered Pairs

Definition 2.1.1 (Ordered Pair). For any sets a and b , the *ordered pair* (a, b) is defined to be $\{\{a\}, \{a, b\}\}$.

Theorem 2.1.2 (Z). For any sets a, b, c, d , we have $(a, b) = (c, d)$ if and only if $a = c$ and $b = d$.

PROOF:

$\langle 1 \rangle 1$. If $(a, b) = (c, d)$ then $a = c$ and $b = d$.

$\langle 2 \rangle 1$. ASSUME: $\{\{a\}, \{a, b\}\} = \{\{c\}, \{c, d\}\}$

$\langle 2 \rangle 2$. $\bigcap \{\{a\}, \{a, b\}\} = \bigcap \{\{c\}, \{c, d\}\}$

$\langle 2 \rangle 3$. $\{a\} = \{c\}$

$\langle 2 \rangle 4$. $a = c$

$\langle 2 \rangle 5$. $\bigcup \{\{a\}, \{a, b\}\} = \bigcup \{\{c\}, \{c, d\}\}$

$\langle 2 \rangle 6$. $\{a, b\} = \{c, d\}$

$\langle 2 \rangle 7$. $b = c$ or $b = d$

$\langle 2 \rangle 8$. $a = d$ or $b = d$

$\langle 2 \rangle 9$. If $b = c$ and $a = d$ then $b = d$

PROOF: By $\langle 2 \rangle 4$.

$\langle 2 \rangle 10$. $b = d$

PROOF: From $\langle 2 \rangle 7$, $\langle 2 \rangle 8$, $\langle 2 \rangle 9$.

$\langle 1 \rangle 2$. If $a = c$ and $b = d$ then $(a, b) = (c, d)$.

PROOF: First-order logic.

□

Definition 2.1.3 (Cartesian Product). The *Cartesian product* of classes \mathbf{A} and \mathbf{B} is the class $\mathbf{A} \times \mathbf{B} := \{(x, y) \mid x \in \mathbf{A}, y \in \mathbf{B}\}$.

Proposition 2.1.4 (Z). For any sets A and B , the class $A \times B$ is a set.

PROOF: It is a subset of $\mathcal{PP}(A \cup B)$. □

Proposition Schema 2.1.5 (Z). *For any classes \mathbf{A} , \mathbf{B} and \mathbf{C} , the following is a theorem:*

$$\mathbf{A} \times (\mathbf{B} \cup \mathbf{C}) = (\mathbf{A} \times \mathbf{B}) \cup (\mathbf{A} \times \mathbf{C})$$

PROOF:

$$\begin{aligned} (x, y) \in \mathbf{A} \times (\mathbf{B} \cup \mathbf{C}) &\Leftrightarrow x \in \mathbf{A} \wedge (y \in \mathbf{B} \vee y \in \mathbf{C}) \\ &\Leftrightarrow (x \in \mathbf{A} \wedge y \in \mathbf{B}) \vee (x \in \mathbf{A} \wedge y \in \mathbf{C}) \\ &\Leftrightarrow (x, y) \in (\mathbf{A} \times \mathbf{B}) \cup (\mathbf{A} \times \mathbf{C}) \quad \square \end{aligned}$$

Proposition Schema 2.1.6 (Z). *For any classes \mathbf{A} and \mathbf{B} , the following is a theorem:*

If $\mathbf{A} \times \mathbf{B} = \mathbf{A} \times \mathbf{C}$ and \mathbf{A} is nonempty then $\mathbf{B} = \mathbf{C}$.

PROOF:

- $\langle 1 \rangle 1$. PICK $a \in \mathbf{A}$
 $\langle 1 \rangle 2$. For all x we have $x \in \mathbf{B}$ iff $x \in \mathbf{C}$.

PROOF:

$$\begin{aligned} x \in \mathbf{B} &\Leftrightarrow (a, x) \in \mathbf{A} \times \mathbf{B} \\ &\Leftrightarrow (a, x) \in \mathbf{A} \times \mathbf{C} \\ &\Leftrightarrow x \in \mathbf{C} \end{aligned}$$

\square

Proposition Schema 2.1.7 (Z). *For any classes \mathbf{A} and \mathbf{B} , the following is a theorem:*

$$\mathbf{A} \times \bigcup \mathbf{B} = \{(a, b) \mid \exists Y \in \mathbf{B}. (a \in \mathbf{A} \wedge b \in Y)\}$$

PROOF:

$$\begin{aligned} (x, y) \in \mathbf{A} \times \bigcup \mathbf{B} &\Leftrightarrow x \in \mathbf{A} \wedge \exists Y \in \mathbf{B}. y \in Y \\ &\Leftrightarrow \exists Y \in \mathbf{B} (x \in \mathbf{A} \wedge y \in Y) \quad \square \end{aligned}$$

2.2 Relations

Definition 2.2.1 (Relation). A *relation* \mathbf{R} between classes \mathbf{A} and \mathbf{B} is a subclass of $\mathbf{A} \times \mathbf{B}$.

A *(binary) relation on \mathbf{A}* is a relation between \mathbf{A} and \mathbf{A} .

We write $x\mathbf{R}y$ for $(x, y) \in \mathbf{R}$.

2.2.1 Identity Functions

Definition 2.2.2 (Identity Function). For any class \mathbf{A} , the *identity function* or *diagonal relation* $\text{id}_{\mathbf{A}}$ on \mathbf{A} is

$$\text{id}_{\mathbf{A}} := \{(x, x) \mid x \in \mathbf{A}\} .$$

2.2.2 Inverses

Definition 2.2.3 (Inverse). The *inverse* of a relation \mathbf{R} between \mathbf{A} and \mathbf{B} is the relation \mathbf{R}^{-1} between \mathbf{B} and \mathbf{A} defined by

$$b\mathbf{R}^{-1}a \Leftrightarrow a\mathbf{R}b .$$

Proposition Schema 2.2.4 (Z). *For any classes \mathbf{A} , \mathbf{B} and \mathbf{R} , the following is a theorem:*

If \mathbf{R} is a relation between \mathbf{A} and \mathbf{B} , we have $(\mathbf{R}^{-1})^{-1} = \mathbf{R}$.

PROOF:

$$\begin{aligned} x(\mathbf{R}^{-1})^{-1}y &\Leftrightarrow y\mathbf{R}^{-1}x \\ &\Leftrightarrow x\mathbf{R}y \end{aligned}$$

□

2.2.3 Composition

Definition 2.2.5 (Composition). Let \mathbf{R} be a relation between \mathbf{A} and \mathbf{B} , and \mathbf{S} be a relation between \mathbf{B} and \mathbf{C} . The *composition* $\mathbf{S} \circ \mathbf{R}$ is the relation between \mathbf{A} and \mathbf{C} defined by

$$a(\mathbf{S} \circ \mathbf{R})c \Leftrightarrow \exists b(a\mathbf{R}b \wedge b\mathbf{S}c) .$$

Proposition Schema 2.2.6 (Z). *For any classes \mathbf{A} , \mathbf{B} , \mathbf{C} , \mathbf{R} and \mathbf{S} , the following is a theorem:*

If \mathbf{R} is a relation between \mathbf{A} and \mathbf{B} , and \mathbf{S} is a relation between \mathbf{B} and \mathbf{C} , then

$$(\mathbf{S} \circ \mathbf{R})^{-1} = \mathbf{R}^{-1} \circ \mathbf{S}^{-1} .$$

PROOF:

$$\begin{aligned} z(\mathbf{S} \circ \mathbf{R})^{-1}x &\Leftrightarrow x(\mathbf{S} \circ \mathbf{R})z \\ &\Leftrightarrow \exists y.(x\mathbf{R}y \wedge y\mathbf{S}z) \\ &\Leftrightarrow \exists y.(y\mathbf{R}^{-1}x \wedge z\mathbf{S}^{-1}y) \\ &\Leftrightarrow z(\mathbf{R}^{-1} \circ \mathbf{S}^{-1})x \end{aligned}$$

□

2.2.4 Properties of Relations

Definition 2.2.7 (Reflexive). Let \mathbf{R} be a binary relation on \mathbf{A} . Then \mathbf{R} is *reflexive* on \mathbf{A} iff $\forall x \in \mathbf{A} . (x, x) \in \mathbf{R}$.

Proposition Schema 2.2.8 (Z). *For any classes \mathbf{A} and \mathbf{R} , the following is a theorem:*

If \mathbf{R} is a reflexive relation on \mathbf{A} then so is \mathbf{R}^{-1} .

PROOF:

⟨1⟩1. LET: $x \in \mathbf{A}$

⟨1⟩2. $x\mathbf{R}x$

PROOF: Since \mathbf{R} is reflexive.

$\langle 1 \rangle 3. x\mathbf{R}^{-1}x$

□

Definition 2.2.9 (Irreflexive). A relation \mathbf{R} is *irreflexive* iff there is no x such that $(x, x) \in \mathbf{R}$.

Definition 2.2.10 (Symmetric). A relation \mathbf{R} is *symmetric* iff, whenever $x\mathbf{R}y$, then $y\mathbf{R}x$.

Definition 2.2.11 (Antisymmetric). A relation \mathbf{R} is *antisymmetric* iff, whenever $x\mathbf{R}y$ and $y\mathbf{R}x$, then $x = y$.

Proposition Schema 2.2.12 (Z). *For any classes \mathbf{A} and \mathbf{R} , the following is a theorem:*

If \mathbf{R} is an antisymmetric relation on \mathbf{A} then so is \mathbf{R}^{-1} .

PROOF:

$\langle 1 \rangle 1.$ ASSUME: $x\mathbf{R}^{-1}y$ and $y\mathbf{R}^{-1}x$

$\langle 1 \rangle 2.$ $y\mathbf{R}x$ and $x\mathbf{R}y$

$\langle 1 \rangle 3.$ $x = y$

PROOF: Since \mathbf{R} is antisymmetric.

□

Definition 2.2.13 (Transitive). A relation \mathbf{R} is *transitive* iff, whenever $x\mathbf{R}y$ and $y\mathbf{R}z$, then $x\mathbf{R}z$.

Proposition Schema 2.2.14 (Z). *For any classes \mathbf{A} , \mathbf{B} and \mathbf{R} , the following is a theorem:*

If \mathbf{R} is a transitive relation between \mathbf{A} and \mathbf{B} then \mathbf{R}^{-1} is transitive.

PROOF:

$\langle 1 \rangle 1.$ ASSUME: $(x, y), (y, z) \in \mathbf{R}^{-1}$

$\langle 1 \rangle 2.$ $(y, x), (z, y) \in \mathbf{R}$

$\langle 1 \rangle 3.$ $(z, x) \in \mathbf{R}$

$\langle 1 \rangle 4.$ $(x, z) \in \mathbf{R}^{-1}$

□

Proposition 2.2.15 (Z). *For any relation R on a set A , there exists a smallest transitive relation on A that includes R .*

PROOF: The relation is $\bigcap \{S \in \mathcal{P}A^2 \mid R \subseteq S, S \text{ is transitive}\}$. □

Definition 2.2.16 (Transitive Closure). For any relation R on a set A , the *transitive closure* of R is the smallest transitive relation that includes R .

Definition 2.2.17 (Minimal). Let \mathbf{R} be a relation on \mathbf{A} . An element $m \in \mathbf{A}$ is *minimal* iff there is no $x \in \mathbf{A}$ such that $x\mathbf{R}m$.

Definition 2.2.18 (Maximal). Let \mathbf{R} be a relation on \mathbf{A} . An element $m \in \mathbf{A}$ is *maximal* iff there is no $x \in \mathbf{A}$ such that $m\mathbf{R}x$.

2.3 n-ary Relations

Definition Schema 2.3.1. For any sets a_1, \dots, a_n , define the *ordered n -tuple* (a_1, \dots, a_n) by

$$(a_1) := a_1$$

$$(a_1, \dots, a_n, a_{n+1}) = ((a_1, \dots, a_n), a_{n+1})$$

Definition Schema 2.3.2. An *n -ary relation on \mathbf{A}* is a class of ordered n -tuples all of whose components are in \mathbf{A} .

2.4 Well Founded Relations

Definition 2.4.1 (Well Founded). A relation \mathbf{R} on a class \mathbf{A} is *well founded* iff:

- for all $a \in \mathbf{A}$, the class $\{x \in \mathbf{A} \mid x\mathbf{R}a\}$ is a set;
- every nonempty subset of \mathbf{A} has an \mathbf{R} -minimal element.

Proposition 2.4.2 (Z). *For any class \mathbf{A} , the relation $\{(x, y) \in \mathbf{A}^2 \mid x \in y\}$ is well founded.*

PROOF:

$\langle 1 \rangle 1$. For all $a \in \mathbf{A}$, the class $\{x \in \mathbf{A} \mid x \in a\}$ is a set.

PROOF: It is a subclass of a .

$\langle 1 \rangle 2$. Every nonempty subset of \mathbf{A} has an \in -minimal element.

$\langle 2 \rangle 1$. LET: C be a nonempty subset of \mathbf{A}

$\langle 2 \rangle 2$. PICK $m \in C$ such that $m \cap C = \emptyset$

PROOF: Axiom of Regularity.

$\langle 2 \rangle 3$. m is \in -minimal in C .

□

Proposition Schema 2.4.3 (Z). *For any classes \mathbf{A} , \mathbf{B} and \mathbf{R} , the following is a theorem:*

Assume \mathbf{R} is a well founded relation on \mathbf{A} and $\mathbf{B} \subseteq \mathbf{A}$ is nonempty. Then \mathbf{B} has an \mathbf{R} -minimal element.

PROOF:

$\langle 1 \rangle 1$. PICK $b \in \mathbf{B}$

$\langle 1 \rangle 2$. LET: $S = \{x \in \mathbf{B} \mid x\mathbf{R}b\}$

PROOF: S is a set because it is a subclass of $\{x \in \mathbf{A} \mid x\mathbf{R}b\}$.

$\langle 1 \rangle 3$. CASE: $S = \emptyset$

PROOF: In this case b is an \mathbf{R} -minimal element of \mathbf{B} .

$\langle 1 \rangle 4$. CASE: $S \neq \emptyset$

PROOF: In this cases S has an \mathbf{R} -minimal element, which is an \mathbf{R} -minimal element of \mathbf{B} .

□

Proposition Schema 2.4.4 (Z). *For any classes \mathbf{A} , \mathbf{B} and \mathbf{R} , the following is a theorem:*

Assume \mathbf{R} is a well founded relation on \mathbf{B} and $\mathbf{A} \subseteq \mathbf{B}$. Then $\mathbf{R} \cap \mathbf{A}^2$ is a well founded relation on \mathbf{A} .

PROOF:

$\langle 1 \rangle 1$. LET: $\mathbf{R}' = \mathbf{R} \cap \mathbf{A}^2$

$\langle 1 \rangle 2$. For all $a \in \mathbf{A}$, the class $\{x \in \mathbf{A} \mid x\mathbf{R}'a\}$ is a set.

PROOF: By Comprehension since it is a subclass of $\{x \in \mathbf{B} \mid x\mathbf{R}a\}$.

$\langle 1 \rangle 3$. Every nonempty subset of \mathbf{A} has an \mathbf{R}' -minimal element.

PROOF: It is a nonempty subset of \mathbf{B} and so has an \mathbf{R} -minimal element, which is also an \mathbf{R}' -minimal element.

□

Theorem Schema 2.4.5 (Transfinite Induction Principle (Z)). *For any classes \mathbf{A} , \mathbf{B} and \mathbf{R} , the following is a theorem:*

Assume \mathbf{R} is a well founded relation on \mathbf{A} and $\mathbf{B} \subseteq \mathbf{A}$. Assume that, for all $t \in \mathbf{A}$,

$$\{x \in \mathbf{A} \mid x\mathbf{R}t\} \subseteq \mathbf{B} \Rightarrow t \in \mathbf{B} .$$

Then $\mathbf{B} = \mathbf{A}$.

PROOF:

$\langle 1 \rangle 1$. ASSUME: for a contradiction $\mathbf{B} \neq \mathbf{A}$

$\langle 1 \rangle 2$. PICK an \mathbf{R} -minimal element m of $\mathbf{A} - \mathbf{B}$.

PROOF: Proposition 2.4.3.

$\langle 1 \rangle 3$. $\{x \in \mathbf{A} \mid x\mathbf{R}m\} \subseteq \mathbf{B}$

PROOF: By minimality of m .

$\langle 1 \rangle 4$. $m \in \mathbf{B}$

$\langle 1 \rangle 5$. Q.E.D.

PROOF: This is a contradiction.

□

Theorem 2.4.6 (Z). *The transitive closure of a well founded relation on a set is well founded.*

PROOF:

$\langle 1 \rangle 1$. LET: R be a well founded relation on the set A .

$\langle 1 \rangle 2$. LET: R^t be the transitive closure of R .

$\langle 1 \rangle 3$. For any $x, y \in A$, if $xR^t y$ then there exists $z \in A$ such that zRy .

PROOF: $\{(x, y) \in A^2 \mid \exists z \in A. zRy\}$ is a transitive relation on A that includes R .

$\langle 1 \rangle 4$. LET: B be a nonempty subset of A .

$\langle 1 \rangle 5$. PICK an R -minimal element b of B .

$\langle 1 \rangle 6$. b is R^t -minimal in B .

PROOF: If there exists x such that $xR^t b$ then there exists z such that zRb by

$\langle 1 \rangle 3$.

□

Definition 2.4.7 (Initial Segment). Let \mathbf{R} be a relation on \mathbf{A} and $a \in \mathbf{A}$. The *initial segment* up to a is

$$\text{seg } a := \{x \in \mathbf{A} \mid x\mathbf{R}a\} .$$

Theorem Schema 2.4.8 (Transfinite Recursion Theorem Schema (ZFC)). *For any classes \mathbf{A} , \mathbf{R} and any property $G[x, y, z]$, there exists a class \mathbf{F} such that, for any class \mathbf{F}' the following is a theorem:*

Assume that \mathbf{R} is a well-founded relation on \mathbf{A} . Assume that, for any f and t , there exists a unique z such that $G[f, t, z]$. Then $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{V}$ such that, for all $t \in \mathbf{A}$, we have $\mathbf{F} \upharpoonright \text{seg } t$ is a set and

$$G[\mathbf{F} \upharpoonright \text{seg } t, t, \mathbf{F}(t)] .$$

If $\mathbf{F}' : \mathbf{A} \rightarrow \mathbf{V}$ satisfies that, for all $t \in \mathbf{A}$, we have $\mathbf{F}' \upharpoonright \text{seg } t$ is a set and $G[\mathbf{F}' \upharpoonright \text{seg } t, t, \mathbf{F}'(t)]$, then $\mathbf{F}' = \mathbf{F}$.

PROOF:

- $\langle 1 \rangle 1$. For B a subset of \mathbf{A} , let us say a function $v : B \rightarrow \mathbf{V}$ is *acceptable* iff, for all $x \in B$, we have $\text{seg } x \subseteq B$ and $G[v \upharpoonright \text{seg } x, x, v(x)]$
- $\langle 1 \rangle 2$. LET: \mathbf{K} be the class of all acceptable functions.
- $\langle 1 \rangle 3$. LET: $\mathbf{F} = \bigcup \mathbf{K}$
- $\langle 1 \rangle 4$. For all $B, C \subseteq \mathbf{A}$, given $v_1 : B \rightarrow \mathbf{V}$ and $v_2 : C \rightarrow \mathbf{V}$ acceptable and $x \in B \cap C$, we have $v_1(x) = v_2(x)$
 - $\langle 2 \rangle 1$. ASSUME: as transfinite induction hypothesis $\forall y \mathbf{R} x. y \in B \cap C \Rightarrow v_1(y) = v_2(y)$
 - $\langle 2 \rangle 2$. $v_1 \upharpoonright \text{seg } x = v_2 \upharpoonright \text{seg } x$
 - $\langle 2 \rangle 3$. $G[v_1 \upharpoonright \text{seg } x, x, v_1(x)]$
 - $\langle 2 \rangle 4$. $G[v_2 \upharpoonright \text{seg } x, x, v_2(x)]$
 - $\langle 2 \rangle 5$. $v_1(x) = v_2(x)$
- $\langle 1 \rangle 5$. \mathbf{F} is a function.
 - $\langle 2 \rangle 1$. ASSUME: $(x, y), (x, z) \in \mathbf{F}$
 - $\langle 2 \rangle 2$. PICK acceptable $v_1 : B \rightarrow \mathbf{V}$ and $v_2 : C \rightarrow \mathbf{V}$ such that $v_1(x) = y$ and $v_2(x) = z$
 - $\langle 2 \rangle 3$. $y = z$
- PROOF: By $\langle 1 \rangle 4$.
- $\langle 1 \rangle 6$. For all $t \in \text{dom } \mathbf{F}$, we have $\mathbf{F} \upharpoonright \text{seg } t$ is a set and $G[\mathbf{F} \upharpoonright \text{seg } t, t, \mathbf{F}(t)]$
 - $\langle 2 \rangle 1$. LET: $t \in \text{dom } \mathbf{F}$
 - $\langle 2 \rangle 2$. PICK an acceptable $v : A \rightarrow \mathbf{V}$ such that $t \in A$
 - $\langle 2 \rangle 3$. For all $y \mathbf{R} x$ we have $v(y) = \mathbf{F}(y)$
 - $\langle 2 \rangle 4$. $\mathbf{F} \upharpoonright \text{seg } x = v \upharpoonright \text{seg } x$
 - $\langle 2 \rangle 5$. $G[v \upharpoonright \text{seg } x, x, v(x)]$
 - $\langle 2 \rangle 6$. $G[\mathbf{F} \upharpoonright \text{seg } x, x, \mathbf{F}(x)]$
- $\langle 1 \rangle 7$. $\text{dom } \mathbf{F} = \mathbf{A}$
 - $\langle 2 \rangle 1$. LET: $x \in \mathbf{A}$
 - $\langle 2 \rangle 2$. ASSUME: as transfinite induction hypothesis $\forall y \mathbf{R} x. y \in \mathbf{A}$
 - $\langle 2 \rangle 3$. ASSUME: for a contradiction $x \notin \text{dom } \mathbf{F}$

$\langle 2 \rangle 4.$ $\mathbf{F} \upharpoonright \text{seg } x$ is a set

PROOF: Axiom of Replacement.

$\langle 2 \rangle 5.$ $\mathbf{F} \upharpoonright \text{seg } x$ is acceptable

$\langle 2 \rangle 6.$ LET: y be the unique object such that $G[\mathbf{F} \upharpoonright \text{seg } x, x, y]$

$\langle 2 \rangle 7.$ $\mathbf{F} \upharpoonright \text{seg } x \cup \{(x, y)\}$ is acceptable

$\langle 2 \rangle 8.$ $x \in \text{dom } \mathbf{F}$

$\langle 2 \rangle 9.$ Q.E.D.

PROOF: This is a contradiction.

$\langle 1 \rangle 8.$ If $\mathbf{F}' : \mathbf{A} \rightarrow \mathbf{V}$ satisfies the theorem, then $\mathbf{F}' = \mathbf{F}$.

$\langle 2 \rangle 1.$ LET: $x \in \mathbf{A}$

PROVE: $\mathbf{F}'(x) = \mathbf{F}(x)$

$\langle 2 \rangle 2.$ ASSUME: as transfinite induction hypothesis $\forall y \mathbf{R}x. \mathbf{F}'(y) = \mathbf{F}(y)$

$\langle 2 \rangle 3.$ $\mathbf{F} \upharpoonright x = \mathbf{F}' \upharpoonright x$

$\langle 2 \rangle 4.$ $G[\mathbf{F} \upharpoonright x, x, \mathbf{F}(x)]$

$\langle 2 \rangle 5.$ $G[\mathbf{F}' \upharpoonright x, x, \mathbf{F}'(x)]$

$\langle 2 \rangle 6.$ $\mathbf{F}(x) = \mathbf{F}'(x)$

□

Chapter 3

Functions

3.1 Functions

Definition 3.1.1 (Function). A *function* from \mathbf{A} to \mathbf{B} is a relation \mathbf{F} between \mathbf{A} and \mathbf{B} such that, for all $x \in \mathbf{A}$, there is only one y such that $x\mathbf{F}y$. We denote this y by $\mathbf{F}(x)$.

A *binary operation* on a class \mathbf{A} is a function $\mathbf{A}^2 \rightarrow \mathbf{A}$.

Definition 3.1.2 (Closed). Let $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{A}$ be a function and $\mathbf{B} \subseteq \mathbf{A}$. Then \mathbf{B} is *closed* under \mathbf{F} iff $\forall x \in \mathbf{B}. \mathbf{F}(x) \in \mathbf{B}$.

Proposition 3.1.3 (Z). *For any class \mathbf{A} , the following is a theorem:*

$$\text{id}_{\mathbf{A}} : \mathbf{A} \rightarrow \mathbf{A}$$

PROOF: For all $x \in \mathbf{A}$, the only y such that $(x, y) \in \text{id}_{\mathbf{A}}$ is $y = x$. \square

Proposition Schema 3.1.4 (Z). *For any classes \mathbf{A} , \mathbf{B} , \mathbf{C} , \mathbf{F} and \mathbf{G} , the following is a theorem:*

Assume $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{G} : \mathbf{B} \rightarrow \mathbf{C}$. Then $\mathbf{G} \circ \mathbf{F} : \mathbf{A} \rightarrow \mathbf{C}$ and, for all $x \in \mathbf{A}$, we have

$$(\mathbf{G} \circ \mathbf{F})(x) = \mathbf{G}(\mathbf{F}(x)) .$$

PROOF:

$\langle 1 \rangle 1. \forall x \in \mathbf{A}. (x, \mathbf{G}(\mathbf{F}(x))) \in \mathbf{G} \circ \mathbf{F}$

PROOF: Because $(x, \mathbf{F}(x)) \in \mathbf{F}$ and $(\mathbf{F}(x), \mathbf{G}(\mathbf{F}(x))) \in \mathbf{G}$.

$\langle 1 \rangle 2. \text{ If } (x, z) \in \mathbf{F} \circ \mathbf{G} \text{ then } z = \mathbf{G}(\mathbf{F}(x))$

$\langle 2 \rangle 1. \text{ PICK } y \in \mathbf{B} \text{ such that } x\mathbf{F}y \text{ and } y\mathbf{G}z$

$\langle 2 \rangle 2. y = \mathbf{F}(x)$

$\langle 2 \rangle 3. z = \mathbf{G}(y)$

$\langle 2 \rangle 4. z = \mathbf{G}(\mathbf{F}(x))$

\square

Proposition 3.1.5 (Z). *For any set A there exists a function $F : \mathcal{P}A - \{\emptyset\} \rightarrow A$ (a choice function for A) such that, for every nonempty $B \subseteq A$, we have $F(B) \in B$.*

PROOF:

$\langle 1 \rangle 1$. LET: A be a set.

$\langle 1 \rangle 2$. LET: $\mathcal{A} = \{\{B\} \times B \mid B \in \mathcal{P}A - \{\emptyset\}\}$

$\langle 1 \rangle 3$. Every member of \mathcal{A} is nonempty.

$\langle 1 \rangle 4$. Any two distinct members of \mathcal{A} are disjoint.

$\langle 1 \rangle 5$. PICK a set C such that, for all $X \in \mathcal{A}$, we have $C \cap X$ is a singleton.

PROOF: Axiom of Choice.

$\langle 1 \rangle 6$. LET: $F = C \cap \bigcup \mathcal{A}$

$\langle 1 \rangle 7$. $F : \mathcal{P}A - \{\emptyset\} \rightarrow A$

$\langle 2 \rangle 1$. F is a function.

$\langle 3 \rangle 1$. LET: $(B, b), (B, b') \in F$

$\langle 3 \rangle 2$. $(B, b), (B, b') \in \{B\} \times B$

PROOF: Since $(B, b), (B, b') \in \bigcup \mathcal{A}$.

$\langle 3 \rangle 3$. $(B, b), (B, b') \in C \cap (\{B\} \times B)$

$\langle 3 \rangle 4$. $(B, b) = (B, b')$

PROOF: From $\langle 1 \rangle 5$.

$\langle 3 \rangle 5$. $b = b'$

$\langle 2 \rangle 2$. $\text{dom } F = \mathcal{P}A - \{\emptyset\}$

PROOF:

$$B \in \text{dom } F$$

$$\Leftrightarrow \exists b. (B, b) \in F$$

$$\Leftrightarrow \exists b. ((B, b) \in \bigcup \mathcal{A} \wedge (B, b) \in C)$$

$$\Leftrightarrow \exists b. \exists B' \in \mathcal{P}A - \{\emptyset\}. ((B, b) \in \{B'\} \times B' \wedge (B, b) \in C)$$

$$\Leftrightarrow B \in \mathcal{P}A - \{\emptyset\} \wedge \exists b \in B. (B, b) \in C$$

$$\Leftrightarrow B \in \mathcal{P}A - \{\emptyset\} \quad (\langle 1 \rangle 5)$$

$\langle 2 \rangle 3$. $\text{ran } F \subseteq A$

$\langle 1 \rangle 8$. For every nonempty $B \subseteq A$ we have $F(B) \in B$

□

Proposition 3.1.6 (Z). *For any relation R between A and B , there exists a function $H : A \rightarrow B$ such that $H \subseteq R$ (i.e. $\forall x \in A. xRH(x)$).*

PROOF:

$\langle 1 \rangle 1$. LET: R be a relation between A and B .

$\langle 1 \rangle 2$. PICK a choice function G for B .

$\langle 1 \rangle 3$. Define $H : A \rightarrow B$ by $H(x) = G(\{y \mid xRy\})$

$\langle 1 \rangle 4$. $H \subseteq R$

□

3.1.1 Injective Functions

Definition 3.1.7 (Injective). A function $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ is *one-to-one*, *injective* or an *injection*, $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$, iff, for all $x, y \in \mathbf{A}$, if $\mathbf{F}(x) = \mathbf{F}(y)$, then $x = y$.

Proposition 3.1.8 (Z). For any class \mathbf{A} , the following is a theorem:

$\text{id}_{\mathbf{A}} : \mathbf{A} \rightarrow \mathbf{A}$ is injective.

PROOF: If $\text{id}_{\mathbf{A}}(x) = \text{id}_{\mathbf{A}}(y)$ then immediately $x = y$. \square

Proposition Schema 3.1.9 (Z). For any classes $\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{F}, \mathbf{G}$, the following is a theorem:

Assume $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{G} : \mathbf{B} \rightarrow \mathbf{C}$. Then $\mathbf{G} \circ \mathbf{F} : \mathbf{A} \rightarrow \mathbf{C}$.

PROOF:

$\langle 1 \rangle 1$. LET: $x, y \in \mathbf{A}$

$\langle 1 \rangle 2$. ASSUME: $(\mathbf{G} \circ \mathbf{F})(x) = (\mathbf{G} \circ \mathbf{F})(y)$

$\langle 1 \rangle 3$. $\mathbf{G}(\mathbf{F}(x)) = \mathbf{G}(\mathbf{F}(y))$

$\langle 1 \rangle 4$. $\mathbf{F}(x) = \mathbf{F}(y)$

PROOF: Since \mathbf{G} is injective.

$\langle 1 \rangle 5$. $x = y$

PROOF: Since \mathbf{F} is injective.

\square

Proposition 3.1.10 (Z). Let $F : A \rightarrow B$ where A is nonempty. There exists $G : B \rightarrow A$ (a left inverse) such that $G \circ F = \text{id}_A$ if and only if F is one-to-one.

PROOF:

$\langle 1 \rangle 1$. If there exists $G : B \rightarrow A$ such that $G \circ F = \text{id}_A$ then F is one-to-one.

$\langle 2 \rangle 1$. ASSUME: $G : B \rightarrow A$ and $G \circ F = I_A$

$\langle 2 \rangle 2$. LET: $x, y \in A$

$\langle 2 \rangle 3$. ASSUME: $F(x) = F(y)$

$\langle 2 \rangle 4$. $x = y$

PROOF: $x = G(F(x)) = G(F(y)) = y$

$\langle 1 \rangle 2$. If F is one-to-one then there exists $G : B \rightarrow A$ such that $G \circ F = I_A$.

$\langle 2 \rangle 1$. ASSUME: F is one-to-one.

$\langle 2 \rangle 2$. PICK $a \in A$

$\langle 2 \rangle 3$. LET: $G : B \rightarrow A$ be the function defined by: $G(b)$ is the (unique) $x \in A$ such that $F(x) = b$ if there exists such an x , $G(b) = a$ otherwise.

$\langle 2 \rangle 4$. For all $x \in A$ we have $G(F(x)) = x$.

\square

3.1.2 Surjective Functions

Definition 3.1.11 (Surjective). Let $F : A \rightarrow B$. We say that F is *surjective*, or maps A *onto* B , and write $F : A \rightarrow B$, iff for all $y \in B$ there exists $x \in A$ such that $F(x) = y$.

Proposition Schema 3.1.12 (Z). For any class \mathbf{A} , the following is a theorem:

$\text{id}_{\mathbf{A}} : \mathbf{A} \rightarrow \mathbf{A}$ is surjective.

PROOF: For any $y \in \mathbf{A}$ we have $\text{id}_{\mathbf{A}}(y) = y$. \square

Proposition Schema 3.1.13 (Z). *For any classes \mathbf{A} , \mathbf{B} , \mathbf{C} , \mathbf{F} and \mathbf{G} , the following is a theorem:*

If $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{G} : \mathbf{B} \rightarrow \mathbf{C}$, then $\mathbf{G} \circ \mathbf{F} : \mathbf{A} \rightarrow \mathbf{C}$.

PROOF:

- $\langle 1 \rangle 1$. LET: $c \in \mathbf{C}$
- $\langle 1 \rangle 2$. PICK $b \in \mathbf{B}$ such that $\mathbf{G}(b) = c$.
- $\langle 1 \rangle 3$. PICK $a \in \mathbf{A}$ such that $\mathbf{F}(a) = b$.
- $\langle 1 \rangle 4$. $(\mathbf{G} \circ \mathbf{F})(a) = c$

\square

Proposition 3.1.14 (Z). *Let $F : A \rightarrow B$. There exists $H : B \rightarrow A$ (a right inverse) such that $F \circ H = \text{id}_B$ if and only if F maps A onto B .*

PROOF:

- $\langle 1 \rangle 1$. If F has a right inverse then F is surjective.
 - $\langle 2 \rangle 1$. ASSUME: F has a right inverse $H : B \rightarrow A$.
 - $\langle 2 \rangle 2$. LET: $y \in B$
 - $\langle 2 \rangle 3$. $F(H(y)) = y$
 - $\langle 2 \rangle 4$. There exists $x \in A$ such that $F(x) = y$
- $\langle 1 \rangle 2$. If F is surjective then F has a right inverse.
 - $\langle 2 \rangle 1$. ASSUME: F is surjective.
 - $\langle 2 \rangle 2$. PICK a function $H : B \rightarrow A$ such that $H \subseteq F^{-1}$

PROOF: Proposition 3.1.6.

- $\langle 2 \rangle 3$. $F \circ H = \text{id}_B$
 - $\langle 3 \rangle 1$. LET: $y \in B$
 - $\langle 3 \rangle 2$. $(y, H(y)) \in F^{-1}$
 - $\langle 3 \rangle 3$. $F(H(y)) = y$

\square

3.1.3 Bijections

Definition 3.1.15 (Bijection). Let $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$. Then \mathbf{F} is *bijective* or a *bijection*, $\mathbf{F} : \mathbf{A} \approx \mathbf{B}$, iff it is injective and surjective.

Proposition Schema 3.1.16 (Z). *For any class \mathbf{A} , the following is a theorem:*
The identity function $\text{id}_{\mathbf{A}} : \mathbf{A} \approx \mathbf{A}$ is a bijection.

PROOF: Proposition 3.1.8 and 3.1.12. \square

Proposition Schema 3.1.17 (Z). *For any classes \mathbf{A} , \mathbf{B} and \mathbf{F} , the following is a theorem:*

If $\mathbf{F} : \mathbf{A} \approx \mathbf{B}$ then $\mathbf{F}^{-1} : \mathbf{B} \approx \mathbf{A}$.

PROOF:

- $\langle 1 \rangle 1$. $\mathbf{F}^{-1} : \mathbf{B} \rightarrow \mathbf{A}$
- $\langle 2 \rangle 1$. LET: $b \in \mathbf{B}$

⟨2⟩2. PICK $a \in \mathbf{A}$ such that $\mathbf{F}(a) = b$.

PROOF: Since \mathbf{F} is surjective.

⟨2⟩3. $(b, a) \in \mathbf{F}^{-1}$

⟨2⟩4. If $(b, a') \in \mathbf{F}^{-1}$ then $a' = a$.

⟨3⟩1. LET: $a' \in \mathbf{A}$ such that $(b, a') \in \mathbf{F}^{-1}$

⟨3⟩2. $\mathbf{F}(a') = \mathbf{F}(a)$

⟨3⟩3. $a' = a$

PROOF: Since \mathbf{F} is injective.

⟨1⟩2. \mathbf{F}^{-1} is injective.

⟨2⟩1. LET: $x, y \in \mathbf{B}$

⟨2⟩2. ASSUME: $\mathbf{F}^{-1}(x) = \mathbf{F}^{-1}(y)$

⟨2⟩3. $x = y$

PROOF: $x = \mathbf{F}(\mathbf{F}^{-1}(x)) = \mathbf{F}(\mathbf{F}^{-1}(y)) = y$.

⟨1⟩3. \mathbf{F}^{-1} is surjective.

PROOF: For all $a \in \mathbf{A}$ we have $\mathbf{F}^{-1}(\mathbf{F}(a)) = a$.

□

Proposition Schema 3.1.18 (Z). *For any classes \mathbf{A} , \mathbf{B} , \mathbf{C} , \mathbf{F} and \mathbf{G} , the following is a theorem:*

If $\mathbf{F} : \mathbf{A} \approx \mathbf{B}$ and $\mathbf{G} : \mathbf{B} \approx \mathbf{C}$ then $\mathbf{G} \circ \mathbf{F} : \mathbf{A} \approx \mathbf{C}$.

PROOF: Propositions 3.1.9 and 3.1.13. □

3.1.4 Restrictions

Definition 3.1.19 (Restriction). Let $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$. Let $\mathbf{C} \subseteq \mathbf{A}$. The *restriction* of \mathbf{F} to \mathbf{C} , denoted $\mathbf{F} \upharpoonright \mathbf{C}$, is the function

$$\begin{aligned} \mathbf{F} \upharpoonright \mathbf{C} : \mathbf{C} &\rightarrow \mathbf{B} \\ (\mathbf{F} \upharpoonright \mathbf{C})(x) &= \mathbf{F}(x) \quad (x \in \mathbf{C}) \end{aligned}$$

3.1.5 Images

Definition 3.1.20 (Image). Let $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{C} \subseteq \mathbf{A}$. The *image* of \mathbf{C} under \mathbf{F} is the class

$$\mathbf{F}(\mathbf{C}) := \{\mathbf{F}(x) \mid x \in \mathbf{C}\}.$$

Proposition Schema 3.1.21 (Z). *For any classes \mathbf{F} , \mathbf{A} and \mathbf{B} , the following is a theorem.*

If $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$, then for any subset $S \subseteq \mathbf{A}$, the class $\mathbf{F}(S)$ is a set.

PROOF: By an Axiom of Replacement. □

Proposition Schema 3.1.22 (Z). *For any classes \mathbf{A} , \mathbf{B} , \mathbf{C} and \mathbf{F} , the following is a theorem:*

Assume $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{C} \subseteq \mathcal{P}\mathbf{A}$. Then

$$\mathbf{F}\left(\bigcup \mathbf{C}\right) = \{y \mid \exists X \in \mathbf{C}. y \in \mathbf{F}(X)\}$$

PROOF:

$$\begin{aligned}
 y \in \mathbf{F}\left(\bigcup \mathbf{C}\right) &\Leftrightarrow \exists x \in \bigcup \mathbf{C}. y = \mathbf{F}(x) \\
 &\Leftrightarrow \exists x. \exists X. X \in \mathbf{C} \wedge x \in X \wedge y = \mathbf{F}(x) \\
 &\Leftrightarrow \exists X \in \mathbf{C}. y \in \mathbf{F}(X) \quad \square
 \end{aligned}$$

Proposition Schema 3.1.23 (Z). *For any classes \mathbf{A} , \mathbf{B} , \mathbf{C} , \mathbf{D} and \mathbf{F} , the following is a theorem:*

Assume $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{C}, \mathbf{D} \subseteq \mathbf{A}$. Then

$$\mathbf{F}(\mathbf{C} \cup \mathbf{D}) = \mathbf{F}(\mathbf{C}) \cup \mathbf{F}(\mathbf{D}) .$$

PROOF:

$$\begin{aligned}
 y \in \mathbf{F}(\mathbf{C} \cup \mathbf{D}) &\Leftrightarrow \exists x \in \mathbf{C} \cup \mathbf{D}. y = \mathbf{F}(x) \\
 &\Leftrightarrow \exists x \in \mathbf{C}. y = \mathbf{F}(x) \vee \exists x \in \mathbf{D}. y = \mathbf{F}(x) \\
 &\Leftrightarrow y \in \mathbf{F}(\mathbf{C}) \cup \mathbf{F}(\mathbf{D}) \quad \square
 \end{aligned}$$

Proposition 3.1.24 (Z). *For any classes \mathbf{F} , \mathbf{A} , \mathbf{B} , \mathbf{C} and \mathbf{D} , the following is a theorem:*

Assume $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{C}, \mathbf{D} \subseteq \mathbf{A}$. Then

$$\mathbf{F}(\mathbf{A} \cap \mathbf{B}) \subseteq \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B}) .$$

Equality holds if \mathbf{F} is injective.

PROOF:

- $\langle 1 \rangle 1. \mathbf{F}(\mathbf{A} \cap \mathbf{B}) \subseteq \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B})$
- $\langle 2 \rangle 1. \text{ LET: } y \in \mathbf{F}(\mathbf{A} \cap \mathbf{B})$
- $\langle 2 \rangle 2. \text{ PICK } x \in \mathbf{A} \cap \mathbf{B} \text{ such that } y = \mathbf{F}(x)$
- $\langle 2 \rangle 3. y \in \mathbf{F}(\mathbf{A})$
- PROOF: Since $x \in \mathbf{A}$.
- $\langle 2 \rangle 4. y \in \mathbf{F}(\mathbf{B})$
- PROOF: Since $x \in \mathbf{B}$.
- $\langle 1 \rangle 2. \text{ If } \mathbf{F} \text{ is injective then } \mathbf{F}(\mathbf{A} \cap \mathbf{B}) = \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B}).$
- $\langle 2 \rangle 1. \text{ ASSUME: } \mathbf{F} \text{ is injective.}$
- $\langle 2 \rangle 2. \text{ LET: } y \in \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B})$
- $\langle 2 \rangle 3. \text{ PICK } x \in \mathbf{A} \text{ such that } y = \mathbf{F}(x)$
- $\langle 2 \rangle 4. \text{ PICK } x' \in \mathbf{B} \text{ such that } y = \mathbf{F}(x')$
- $\langle 2 \rangle 5. x = x'$
- PROOF: $\langle 2 \rangle 1$
- $\langle 2 \rangle 6. x \in \mathbf{A} \cap \mathbf{B}$
- $\langle 2 \rangle 7. y \in \mathbf{F}(\mathbf{A} \cap \mathbf{B})$

□

Proposition Schema 3.1.25 (Z). *For any classes \mathbf{F} , \mathbf{A} , \mathbf{B} , and \mathbf{C} , the following is a theorem:*

Let $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{C} \subseteq \mathcal{P}\mathbf{A}$. Then

$$\mathbf{F}\left(\bigcap \mathbf{C}\right) \subseteq \bigcap \{\mathbf{F}(X) \mid X \in \mathbf{A}\} .$$

Equality holds if \mathbf{F} is injective and \mathbf{A} is nonempty.

PROOF:

- $\langle 1 \rangle 1.$ $\mathbf{F}(\bigcap \mathbf{A}) \subseteq \bigcap \{\mathbf{F}(X) \mid X \in \mathbf{A}\}$
 - $\langle 2 \rangle 1.$ LET: $y \in \mathbf{F}(\bigcap \mathbf{A})$
 - $\langle 2 \rangle 2.$ PICK $x \in \bigcap \mathbf{A}$ such that $y = \mathbf{F}(x)$
 - $\langle 2 \rangle 3.$ LET: $X \in \mathbf{A}$
 - $\langle 2 \rangle 4.$ $x \in X$
 - $\langle 2 \rangle 5.$ $y \in \mathbf{F}(X)$
- $\langle 1 \rangle 2.$ If \mathbf{F} is injective then $\mathbf{F}(\bigcap \mathbf{A}) = \bigcap \{\mathbf{F}(X) \mid X \in \mathbf{A}\}$
 - $\langle 2 \rangle 1.$ ASSUME: \mathbf{F} is injective.
 - $\langle 2 \rangle 2.$ ASSUME: \mathbf{A} is nonempty.
 - $\langle 2 \rangle 3.$ LET: $y \in \bigcap \{\mathbf{F}(X) \mid X \in \mathbf{A}\}$
 - $\langle 2 \rangle 4.$ PICK $X_0 \in \mathbf{A}$
 - $\langle 2 \rangle 5.$ PICK $x \in X_0$ such that $(x, y) \in \mathbf{F}$
 - $\langle 2 \rangle 6.$ $x \in \bigcap \mathbf{A}$
 - $\langle 3 \rangle 1.$ LET: $X \in \mathbf{A}$
 - $\langle 3 \rangle 2.$ PICK $x' \in X$ such that $(x', y) \in \mathbf{F}$.
 - $\langle 3 \rangle 3.$ $x = x'$
 - PROOF: $\langle 2 \rangle 1$
 - $\langle 3 \rangle 4.$ $x \in X$
 - $\langle 2 \rangle 7.$ $y \in \mathbf{F}(\bigcap \mathbf{A})$

□

Proposition 3.1.26 (Z). For any classes $\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D}$ and \mathbf{F} , the following is a theorem:

Assume $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{C}, \mathbf{D} \subseteq \mathbf{A}$. Then

$$\mathbf{F}(\mathbf{C}) - \mathbf{F}(\mathbf{D}) \subseteq \mathbf{F}(\mathbf{C} - \mathbf{D}) .$$

Equality holds if \mathbf{F} is injective.

PROOF:

- $\langle 1 \rangle 1.$ $\mathbf{F}(\mathbf{C}) - \mathbf{F}(\mathbf{D}) \subseteq \mathbf{F}(\mathbf{A} - \mathbf{B})$
 - $\langle 2 \rangle 1.$ LET: $y \in \mathbf{F}(\mathbf{A}) - \mathbf{F}(\mathbf{B})$
 - $\langle 2 \rangle 2.$ PICK $x \in \mathbf{A}$ such that $y = \mathbf{F}(x)$
 - $\langle 2 \rangle 3.$ $x \notin \mathbf{B}$
 - $\langle 2 \rangle 4.$ $x \in \mathbf{A} - \mathbf{B}$
 - $\langle 2 \rangle 5.$ $y \in \mathbf{F}(\mathbf{A} - \mathbf{B})$
- $\langle 1 \rangle 2.$ If \mathbf{F} is injective then $\mathbf{F}(\mathbf{A}) - \mathbf{F}(\mathbf{B}) = \mathbf{F}(\mathbf{A} - \mathbf{B})$
 - $\langle 2 \rangle 1.$ ASSUME: \mathbf{F} is injective.
 - $\langle 2 \rangle 2.$ LET: $y \in \mathbf{F}(\mathbf{A} - \mathbf{B})$
 - $\langle 2 \rangle 3.$ PICK $x \in \mathbf{A} - \mathbf{B}$ such that $y = \mathbf{F}(x)$
 - $\langle 2 \rangle 4.$ $y \in \mathbf{F}(\mathbf{A})$
 - $\langle 2 \rangle 5.$ $y \notin \mathbf{F}(\mathbf{B})$

- ⟨3⟩1. ASSUME: for a contradiction $y \in \mathbf{F}(\mathbf{B})$
 ⟨3⟩2. PICK $x' \in \mathbf{B}$ such that $y = \mathbf{F}(x')$
 ⟨3⟩3. $x = x'$
 PROOF: ⟨2⟩1
 ⟨3⟩4. $x \in \mathbf{B}$
 ⟨3⟩5. Q.E.D.
 PROOF: This contradicts ⟨2⟩3.

□

3.1.6 Inverse Images

Definition 3.1.27 (Inverse Image). Let $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{C} \subseteq \mathbf{B}$. Then the *inverse image* of \mathbf{C} under \mathbf{F} is

$$\mathbf{F}^{-1}(\mathbf{C}) = \{x \in \mathbf{A} \mid \mathbf{F}(x) \in \mathbf{C}\} .$$

Proposition Schema 3.1.28 (Z). For any classes \mathbf{A} , \mathbf{B} , \mathbf{C} and \mathbf{F} , the following is a theorem:

Assume $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{C} \subseteq \mathcal{PB}$. Then

$$\mathbf{F}^{-1}\left(\bigcap \mathbf{C}\right) = \bigcap \{\mathbf{F}^{-1}(X) \mid X \in \mathbf{C}\} .$$

PROOF:

$$\begin{aligned}
 x \in \mathbf{F}^{-1}\left(\bigcap \mathbf{C}\right) &\Leftrightarrow \mathbf{F}(x) \in \bigcap \mathbf{C} \\
 &\Leftrightarrow \forall X \in \mathbf{C}. \mathbf{F}(x) \in X \\
 &\Leftrightarrow \forall X \in \mathbf{C}. x \in \mathbf{F}^{-1}(X) \quad \square
 \end{aligned}$$

Proposition Schema 3.1.29 (Z). For any classes \mathbf{A} , \mathbf{B} , \mathbf{C} , \mathbf{D} and \mathbf{F} , the following is a theorem:

Assume $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and $\mathbf{C}, \mathbf{D} \subseteq \mathbf{B}$. Then

$$\mathbf{F}^{-1}(\mathbf{C} - \mathbf{D}) = \mathbf{F}^{-1}(\mathbf{C}) - \mathbf{F}^{-1}(\mathbf{D}) .$$

PROOF:

$$\begin{aligned}
 x \in \mathbf{F}^{-1}(\mathbf{C} - \mathbf{D}) &\Leftrightarrow \mathbf{F}(x) \in \mathbf{C} - \mathbf{D} \\
 &\Leftrightarrow \mathbf{F}(x) \in \mathbf{C} \wedge \mathbf{F}(x) \notin \mathbf{D} \\
 &\Leftrightarrow x \in \mathbf{F}^{-1}(\mathbf{C}) \wedge x \notin \mathbf{F}^{-1}(\mathbf{D}) \\
 &\Leftrightarrow x \in \mathbf{F}^{-1}(\mathbf{C}) - \mathbf{F}^{-1}(\mathbf{D}) \quad \square
 \end{aligned}$$

3.1.7 Function Sets

Proposition 3.1.30 (ZFC). For any classes \mathbf{B} and \mathbf{F} , the following is a theorem:

Let A be a set. If $\mathbf{F} : A \rightarrow \mathbf{B}$ then \mathbf{F} is a set.

PROOF: By an Axiom of Replacement, we have $R = \{\mathbf{F}(x) \mid x \in A\}$ is a set. Hence \mathbf{F} is a set since $\mathbf{F} \subseteq A \times R$. □

Definition 3.1.31 (Dependent Product Class). Let I be a set and let $\mathbf{H}(i)$ be a class for all $i \in I$. We write $\prod_{i \in I} \mathbf{H}(i)$ for the class of all functions $f : I \rightarrow \bigcup_{i \in I} \mathbf{H}(i)$ such that $\forall i \in I. f(i) \in \mathbf{H}(i)$.

We write \mathbf{B}^I for $\prod_{i \in I} \mathbf{B}$ where \mathbf{B} does not depend on I .

Proposition Schema 3.1.32 (ZFC). Let I be a set. Let $H(i)$ be a set for every $i \in I$. Then $\prod_{i \in I} H(i)$ is a set.

PROOF:

$\langle 1 \rangle 1$. $\{\mathbf{H}(i) \mid i \in I\}$ is a set.

PROOF: By an Axiom of Replacement.

$\langle 1 \rangle 2$. $\bigcup_{i \in I} \mathbf{H}(i)$ is a set.

$\langle 1 \rangle 3$. $\prod_{i \in I} \mathbf{H}(i)$ is a set.

PROOF: It is a subset of $\mathcal{P}(I \times \bigcup_{i \in I} \mathbf{H}(i))$.

□

Proposition 3.1.33 (Z). Let I be a set. Let $H(i)$ be a set for all $i \in I$. If $\forall i \in I. H(i) \neq \emptyset$ then $\prod_{i \in I} H(i) \neq \emptyset$.

PROOF:

$\langle 1 \rangle 1$. ASSUME: $\forall i \in I. H(i) \neq \emptyset$

$\langle 1 \rangle 2$. LET: $R = \{(i, x) \mid i \in I, x \in H(i)\}$

$\langle 1 \rangle 3$. PICK a function $f : I \rightarrow \bigcup_{i \in I} H(i)$ such that $f \subseteq R$

PROOF: Proposition 3.1.6.

$\langle 1 \rangle 4$. $f \in \prod_{i \in I} H(i)$

□

3.2 Equinumerosity

Definition 3.2.1 (Equinumerous). Sets A and B are *equinumerous*, $A \approx B$, iff there exists a bijection between A and B .

3.3 Domination

Definition 3.3.1 (Dominate). A set A is *dominated* by a set B , $A \preccurlyeq B$, iff there exists an injection $A \rightarrow B$.

Proposition 3.3.2 (Z). Given sets A and B , if $A \neq \emptyset$ or $B = \emptyset$, then we have $A \preccurlyeq B$ iff there exists a surjective function $B \rightarrow A$.

PROOF:

$\langle 1 \rangle 1$. If $A \preccurlyeq B$ and $A \neq \emptyset$ then there exists a surjective function $B \rightarrow A$.

$\langle 2 \rangle 1$. ASSUME: $f : A \rightarrow B$ be injective.

$\langle 2 \rangle 2$. PICK $a \in A$

$\langle 2 \rangle 3$. LET: $g : B \rightarrow A$ be the function defined by $g(b) = f^{-1}(b)$ if $b \in \text{ran } f$,
and $g(b) = a$ otherwise.

- $\langle 2 \rangle 4$. g is surjective.
- $\langle 1 \rangle 2$. If there exists a surjective function $B \rightarrow A$ then $A \preceq B$.
- $\langle 2 \rangle 1$. ASSUME: there exists a surjective function $g : B \rightarrow A$
- $\langle 2 \rangle 2$. $\forall a \in A. \exists b \in B. g(b) = a$
- $\langle 2 \rangle 3$. Choose a function $f : A \rightarrow B$ such that $\forall a \in A. g(f(a)) = a$
- $\langle 2 \rangle 4$. f is injective.

□

Chapter 4

Equivalence Relations

Definition 4.0.1 (Equivalence Relation). An *equivalence relation* on a class \mathbf{A} is a binary relation on \mathbf{A} that is reflexive, symmetric and transitive.

Proposition 4.0.2 (Z). *Equinumerosity is an equivalence relation on the class of all sets.*

PROOF: Propositions 3.1.16, 3.1.17, 3.1.18. \square

Definition 4.0.3 (Respects). Let \mathbf{R} be an equivalence relation on \mathbf{A} and $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$. Then \mathbf{F} *respects* \mathbf{A} iff, whenever $(x, y) \in \mathbf{R}$, then $\mathbf{F}(x) = \mathbf{F}(y)$.

Definition 4.0.4 (Equivalence Class). Let \mathbf{R} be an equivalence relation on \mathbf{A} and $a \in \mathbf{A}$. The *equivalence class* of a modulo \mathbf{R} is

$$[a]_{\mathbf{R}} := \{x \mid a\mathbf{R}x\} .$$

Proposition Schema 4.0.5 (Z). *For any classes \mathbf{A} and \mathbf{R} , the following is a theorem.*

Assume \mathbf{R} be an equivalence relation on \mathbf{A} . Let $a, b \in \mathbf{A}$. Then $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$ if and only if $a\mathbf{R}b$.

PROOF:

$\langle 1 \rangle 1$. If $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$ then $a\mathbf{R}b$.

$\langle 2 \rangle 1$. ASSUME: $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$

$\langle 2 \rangle 2$. $b\mathbf{R}b$

PROOF: Reflexivity

$\langle 2 \rangle 3$. $b \in [b]_{\mathbf{R}}$

$\langle 2 \rangle 4$. $b \in [a]_{\mathbf{R}}$

$\langle 2 \rangle 5$. $a\mathbf{R}b$

$\langle 1 \rangle 2$. If $a\mathbf{R}b$ then $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$.

$\langle 2 \rangle 1$. For all $x, y \in \mathbf{A}$, if $x\mathbf{R}y$ then $[y]_{\mathbf{R}} \subseteq [x]_{\mathbf{R}}$

$\langle 3 \rangle 1$. LET: $x, y \in \mathbf{A}$

$\langle 3 \rangle 2$. ASSUME: $x\mathbf{R}y$

$\langle 3 \rangle 3$. LET: $t \in [y]_{\mathbf{R}}$

$\langle 3 \rangle 4$. $y \mathbf{R} t$

$\langle 3 \rangle 5$. $x \mathbf{R} t$

PROOF: Transitivity, $\langle 3 \rangle 2$, $\langle 3 \rangle 4$.

$\langle 3 \rangle 6$. $t \in [x]_{\mathbf{R}}$

$\langle 2 \rangle 2$. ASSUME: $a \mathbf{R} b$

$\langle 2 \rangle 3$. $[b]_{\mathbf{R}} \subseteq [a]_{\mathbf{R}}$

PROOF: $\langle 2 \rangle 1$, $\langle 2 \rangle 2$.

$\langle 2 \rangle 4$. $b \mathbf{R} a$

PROOF: Symmetry, $\langle 2 \rangle 2$.

$\langle 2 \rangle 5$. $[a]_{\mathbf{R}} \subseteq [b]_{\mathbf{R}}$

PROOF: $\langle 2 \rangle 1$, $\langle 2 \rangle 4$.

$\langle 2 \rangle 6$. $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$

PROOF: $\langle 2 \rangle 3$, $\langle 2 \rangle 5$.

□

Definition 4.0.6 (Partition). A *partition* Π of a set A is a set of nonempty subsets of A that is disjoint and exhaustive, i.e.

1. no two different sets in Π have any common elements, and
2. each element of A is in some set in Π .

Definition 4.0.7. Let R be an equivalence relation on a set A . The *quotient set* A/R is the set of all equivalence classes.

Theorem 4.0.8 (Z). Let A be a set and \mathbf{B} a class. Let R be an equivalence relation on A and $F : A \rightarrow \mathbf{B}$. Then F respects R if and only if there exists $\hat{F} : A/R \rightarrow \mathbf{B}$ such that

$$\forall a \in A. \hat{F}([a]_R) = F(a) \text{ .}$$

In this case, \hat{F} is unique.

PROOF:

$\langle 1 \rangle 1$. If F respects R then there exists $\hat{F} : A/R \rightarrow \mathbf{B}$ such that $\forall a \in A. \hat{F}([a]_R) = F(a)$.

$\langle 2 \rangle 1$. ASSUME: F respects R .

$\langle 2 \rangle 2$. LET: $\hat{F} = \{([a]_R, F(a)) \mid a \in A\}$

$\langle 2 \rangle 3$. \hat{F} is a function.

$\langle 3 \rangle 1$. ASSUME: $a, a' \in A$ and $[a]_R = [a']_R$

PROVE: $F(a) = F(a')$

$\langle 3 \rangle 2$. $(a, a') \in R$

PROOF: Proposition 4.0.5.

$\langle 3 \rangle 3$. $F(a) = F(a')$

PROOF: $\langle 2 \rangle 1$

$\langle 2 \rangle 4$. $\text{dom } \hat{F} = A/R$

$\langle 2 \rangle 5$. $\text{ran } \hat{F} \subseteq \mathbf{B}$

- $\langle 2 \rangle 6. \forall a \in A. \hat{F}([a]_R) = F(a)$
 $\langle 1 \rangle 2.$ If there exists $\hat{F} : A/R \rightarrow \mathbf{B}$ such that $\forall a \in A. \hat{F}([a]_R) = F(a)$ then F respects R .
 $\langle 2 \rangle 1.$ ASSUME: $\hat{F} : A/R \rightarrow \mathbf{B}$ and $\forall a \in A. \hat{F}([a]_R) = F(a)$
 $\langle 2 \rangle 2.$ LET: $a, a' \in A$
 $\langle 2 \rangle 3.$ ASSUME: $(a, a') \in R$
 $\langle 2 \rangle 4.$ $[a]_R = [a']_R$
 PROOF: Proposition 4.0.5.
 $\langle 2 \rangle 5.$ $F(a) = F(a')$
 PROOF: $\langle 2 \rangle 1$
 $\langle 1 \rangle 3.$ If $G, H : A/R \rightarrow \mathbf{B}$ and $\forall a \in A. G([a]_R) = H([a]_R)$ then $G = H$.
 \square

Proposition 4.0.9 (Z). *Let R be an equivalence relation on a set A . Then A/R is a partition of A .*

PROOF:

- $\langle 1 \rangle 1.$ Every member of A/R is nonempty.
 PROOF: Since $a \in [a]_R$ by reflexivity.
 $\langle 1 \rangle 2.$ No two different sets in A/R have any common elements.
 $\langle 2 \rangle 1.$ LET: $[a]_R, [b]_R \in A/R$
 $\langle 2 \rangle 2.$ LET: $c \in [a]_R \cap [b]_R$
 PROVE: $[a]_R = [b]_R$
 $\langle 2 \rangle 3.$ aRc
 PROOF: $\langle 2 \rangle 2$
 $\langle 2 \rangle 4.$ bRc
 PROOF: $\langle 2 \rangle 2$
 $\langle 2 \rangle 5.$ cRb
 PROOF: Symmetry, $\langle 2 \rangle 4$
 $\langle 2 \rangle 6.$ aRb
 PROOF: Transitivity, $\langle 2 \rangle 3$, $\langle 2 \rangle 5$
 $\langle 2 \rangle 7.$ $[a]_R = [b]_R$
 PROOF: Proposition 4.0.5, $\langle 2 \rangle 6$
 $\langle 1 \rangle 3.$ Each element of A is in some set in A/R .
 PROOF: Since $a \in [a]_R$ by reflexivity.
 \square

Proposition 4.0.10 (Z). *For any partition P of a set A , there exists a unique equivalence relation R on A such that $A/R = P$, namely xRy iff $\exists X \in P(x \in X \wedge y \in X)$.*

PROOF: Easy. \square

Definition 4.0.11 (Natural Map). Let A be a set and R an equivalence relation on A . The *natural map* $A \rightarrow A/R$ is the function that maps $a \in A$ to $[a]_R$.

Chapter 5

Ordering Relations

5.1 Partial Orders

Definition 5.1.1 (Partial Ordering). Let \mathbf{A} be a class. A *partial ordering* on \mathbf{A} is a relation \mathbf{R} on \mathbf{A} that is reflexive, antisymmetric and transitive.

We often write \leq for a partial ordering, and then write $x < y$ for $x \leq y \wedge x \neq y$.

Proposition Schema 5.1.2 (Z). *For any classes \mathbf{A} and \mathbf{R} , the following is a theorem:*

If \mathbf{R} is a partial order on \mathbf{A} then so is \mathbf{R}^{-1} .

PROOF:

$\langle 1 \rangle 1.$ \mathbf{R}^{-1} is reflexive.

PROOF: Proposition 2.2.8.

$\langle 1 \rangle 2.$ \mathbf{R}^{-1} is antisymmetric.

PROOF: Proposition 2.2.12.

$\langle 1 \rangle 3.$ \mathbf{R}^{-1} is transitive.

$\langle 2 \rangle 1.$ ASSUME: $x\mathbf{R}^{-1}y$ and $y\mathbf{R}^{-1}z$

$\langle 2 \rangle 2.$ $y\mathbf{R}x$ and $z\mathbf{R}y$

$\langle 2 \rangle 3.$ $z\mathbf{R}x$

PROOF: Since \mathbf{R} is transitive.

$\langle 2 \rangle 4.$ $x\mathbf{R}^{-1}z$

□

Proposition Schema 5.1.3 (Z). *For any classes \mathbf{A} , \mathbf{B} , \mathbf{F} and \mathbf{R} , the following is a theorem:*

Assume \mathbf{R} is a partial order on \mathbf{B} and $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ is injective. Define \mathbf{S} on \mathbf{A} by $x\mathbf{S}y$ iff $\mathbf{F}(x)\mathbf{R}\mathbf{F}(y)$. Then \mathbf{S} is a partial order on \mathbf{A} .

PROOF:

$\langle 1 \rangle 1.$ \mathbf{S} is reflexive.

PROOF: For any $x \in \mathbf{A}$ we have $\mathbf{F}(x)\mathbf{R}\mathbf{F}(x)$.

- ⟨1⟩2. **S** is antisymmetric.
 ⟨2⟩1. LET: $x, y \in \mathbf{A}$
 ⟨2⟩2. ASSUME: $x\mathbf{S}y$ and $y\mathbf{S}x$
 ⟨2⟩3. $\mathbf{F}(x)\mathbf{R}\mathbf{F}(y)$ and $\mathbf{F}(y)\mathbf{R}\mathbf{F}(x)$
 ⟨2⟩4. $\mathbf{F}(x) = \mathbf{F}(y)$
 PROOF: **R** is antisymmetric.
 ⟨2⟩5. $x = y$
 ⟨1⟩3. **S** is transitive.
 □

Corollary Schema 5.1.3.1 (Z). *For any classes **A**, **B** and **R**, the following is a theorem:*

*Assume **R** be a partial order on **A** and $\mathbf{B} \subseteq \mathbf{A}$. Then $\mathbf{R} \cap \mathbf{B}^2$ is a partial order on **B**.*

Definition 5.1.4 (Partially Ordered Set). A *partially ordered set* or *poset* is a pair (A, \leq) where A is a set and \leq is a partial ordering on A . We often write just A for (A, \leq) .

If (A, \leq) is a poset and $B \subseteq A$ we write just B for the poset $(B, \leq \cap B^2)$.

Definition 5.1.5 (Strictly Monotone). Let $(A, <_A)$ and $(B, <_B)$ be posets. A function $f : A \rightarrow B$ is *strictly monotone* iff, whenever $x <_A y$, then $f(x) <_B f(y)$.

Definition 5.1.6 (Least). Let \leq be a partial order on **A**. An element $m \in \mathbf{A}$ is *least* iff for all $x \in \mathbf{A}$ we have $m \leq x$.

Proposition 5.1.7 (Z). *A partial order has at most one least element.*

PROOF: If m and m' are least then $m \leq m'$ and $m' \leq m$, so $m = m'$. □

Definition 5.1.8 (Greatest). Let \leq be a partial order on **A**. An element $m \in \mathbf{A}$ is *greatest* iff for all $x \in \mathbf{A}$ we have $x \leq m$.

Proposition 5.1.9 (Z). *A poset has at most one greatest element.*

PROOF: If m and m' are greatest then $m \leq m'$ and $m' \leq m$, so $m = m'$. □

Definition 5.1.10 (Upper Bound). Let \leq be a partial ordering on **A** and $\mathbf{B} \subseteq \mathbf{A}$. Let $u \in \mathbf{A}$. Then u is an *upper bound* for **B** iff $\forall x \in \mathbf{B}. x \leq u$.

Definition 5.1.11 (Lower Bound). Let \leq be a partial ordering on **A** and $\mathbf{B} \subseteq \mathbf{A}$. Let $l \in \mathbf{A}$. Then l is a *lower bound* for **B** iff $\forall x \in \mathbf{B}. l \leq x$.

Definition 5.1.12 (Bounded Above). Let \leq be a partial ordering on **A** and $\mathbf{B} \subseteq \mathbf{A}$. Then **B** is *bounded above* iff it has an upper bound.

Definition 5.1.13 (Bounded Below). Let \leq be a partial ordering on **A** and $\mathbf{B} \subseteq \mathbf{A}$. Then **B** is *bounded below* iff it has a lower bound.

Definition 5.1.14 (Least Upper Bound). Let \leq be a partial ordering on \mathbf{A} and $\mathbf{B} \subseteq \mathbf{A}$. Let $s \in \mathbf{A}$. Then s is the *least upper bound* or *supremum* of \mathbf{B} iff s is an upper bound for \mathbf{B} and, for every upper bound u for \mathbf{B} , we have $s \leq u$.

Definition 5.1.15 (Greatest Lower Bound). Let \leq be a partial ordering on \mathbf{A} and $\mathbf{B} \subseteq \mathbf{A}$. Let $i \in \mathbf{A}$. Then i is the *greatest lower bound* or *infimum* of \mathbf{B} iff i is a lower bound for \mathbf{B} and, for every lower bound l for \mathbf{B} , we have $i \leq l$.

Definition 5.1.16 (Complete). A partial order is *complete* iff every nonempty subset bounded above has a supremum, and every nonempty subset bounded below has an infimum.

Definition 5.1.17 (Order Isomorphism). Let A and B be posets. An *order isomorphism* between A and B , $f : A \cong B$, is a bijection $f : A \approx B$ such that f and f^{-1} are monotone.

Theorem 5.1.18 (Knaster Fixed-Point Theorem (Z)). *Let A be a complete poset with a greatest and least element. Let $\phi : A \rightarrow A$ be monotone. Then there exists $a \in A$ such that $\phi(a) = a$.*

PROOF:

$\langle 1 \rangle 1$. LET: $B = \{x \in A \mid x \leq \phi(x)\}$

$\langle 1 \rangle 2$. LET: $a = \sup B$

PROOF: B is nonempty because the least element of A is in B , and it is bounded above by the greatest element of A .

$\langle 1 \rangle 3$. For all $b \in B$ we have $b \leq \phi(a)$

$\langle 2 \rangle 1$. LET: $b \in B$

$\langle 2 \rangle 2$. $b \leq \phi(b)$

$\langle 2 \rangle 3$. $b \leq a$

$\langle 2 \rangle 4$. $\phi(b) \leq \phi(a)$

$\langle 2 \rangle 5$. $b \leq \phi(a)$

$\langle 1 \rangle 4$. $a \leq \phi(a)$

$\langle 1 \rangle 5$. $\phi(a) \leq \phi(\phi(a))$

$\langle 1 \rangle 6$. $\phi(a) \in B$

$\langle 1 \rangle 7$. $\phi(a) \leq a$

$\langle 1 \rangle 8$. $\phi(a) = a$

□

Definition 5.1.19 (Dense). Let \leq be a partial order on \mathbf{A} and $\mathbf{B} \subseteq \mathbf{A}$. Then \mathbf{B} is *dense* iff, for all $x, y \in \mathbf{A}$, if $x < y$ then there exists $z \in \mathbf{B}$ such that $x < z < y$.

Proposition 5.1.20 (Z). *Let A be a complete poset with no least element. Let $B \subseteq A$ be dense. Let $\theta : A \rightarrow A$ be a monotone map that is the identity on B . Then $\theta = \text{id}_A$.*

PROOF:

$\langle 1 \rangle 1$. LET: $a \in A$

PROVE: $\theta(a) = a$

- ⟨1⟩2. LET: $S(a) = \{b \in B \mid b < a\}$
- ⟨1⟩3. $S(a)$ is nonempty and bounded above.
 - ⟨2⟩1. $S(a)$ is nonempty.
 - ⟨3⟩1. PICK $a_1 < a$
 - PROOF: Since a is not least.
 - ⟨3⟩2. There exists $b \in B$ such that $a_1 < b < a$.
 - ⟨2⟩2. $S(a)$ is bounded above by a .
- ⟨1⟩4. $\sup S(a) \leq a$
- ⟨1⟩5. $\sup S(a) = a$
 - ⟨2⟩1. ASSUME: for a contradiction $\sup S(a) < a$
 - ⟨2⟩2. PICK $b \in B$ such that $\sup S(a) < b < a$
 - ⟨2⟩3. $b \in S(a)$
 - ⟨2⟩4. Q.E.D.
- PROOF: This contradicts the fact that $\sup S(a) < b$.
- ⟨1⟩6. For all $b \in S(a)$ we have $b \leq \theta(a)$
 - ⟨2⟩1. LET: $b \in S(a)$
 - ⟨2⟩2. $b < a$
 - ⟨2⟩3. $\theta(b) \leq \theta(a)$
 - ⟨2⟩4. $b \leq \theta(a)$
- PROOF: $\theta(b) = b$
- ⟨1⟩7. $a \leq \theta(a)$
 - PROOF: Since $a = \sup S(a)$ and $\theta(a)$ is an upper bound for $S(a)$.
- ⟨1⟩8. $a \not\leq \theta(a)$
 - ⟨2⟩1. ASSUME: for a contradiction $a < \theta(a)$.
 - ⟨2⟩2. PICK $b \in B$ such that $a < b < \theta(a)$
 - ⟨2⟩3. $\theta(a) \leq \theta(b) = b$
 - ⟨2⟩4. Q.E.D.
- PROOF: This contradicts the fact that $b < \theta(a)$.
- ⟨1⟩9. $\theta(a) = a$

□

Theorem 5.1.21 (Z). *Let A and P be complete posets with no least or greatest element. Let B be dense in A and Q be dense in P . Every order isomorphism $\phi : B \cong Q$ extends uniquely to an order isomorphism $A \cong P$.*

PROOF:

- ⟨1⟩1. For $a \in A$, let $S(a) = \{b \in B \mid b < a\}$.
- ⟨1⟩2. Define $\bar{\phi} : A \rightarrow P$ by $\bar{\phi}(a) = \sup \phi(S(a))$.
 - ⟨2⟩1. $\phi(S(a))$ is nonempty.
 - ⟨3⟩1. PICK $a_1 < a$
 - PROOF: Since a is not least.
 - ⟨3⟩2. PICK $b \in B$ such that $a_1 < b < a$.
 - ⟨3⟩3. $\phi(b) \in \phi(S(a))$
 - ⟨2⟩2. $\phi(S(a))$ is bounded above.
 - ⟨3⟩1. PICK $a_2 > a$
 - PROOF: Since a is not greatest.
 - ⟨3⟩2. PICK $b \in B$ such that $a < b < a_2$

- (3)3. $\phi(b)$ is an upper bound for $\phi(S(a))$.
 (1)3. $\bar{\phi}$ is monotone.
 PROOF: If $a \leq a'$ then $S(a) \subseteq S(a')$ and so $\bar{\phi}(a) \leq \bar{\phi}(a')$.
 (1)4. $\bar{\phi}$ extends ϕ .
 (2)1. LET: $b \in B$
 PROVE: $\phi(b) = \sup \phi(S(b))$
 (2)2. $\phi(b)$ is an upper bound for $\phi(S(b))$
 (2)3. LET: u be any upper bound for $\phi(S(b))$
 PROVE: $\phi(b) \leq u$
 (2)4. ASSUME: for a contradiction $u < \phi(b)$
 (2)5. PICK $q \in Q$ such that $u < q < \phi(b)$
 (2)6. PICK $b' \in B$ such that $\phi(b') = q$
 (2)7. $b' < b$
 (2)8. $b' \in S(b)$
 (2)9. $q = \phi(b') \leq u$
 (2)10. Q.E.D.
 PROOF: This is a contradiction.
 (1)5. LET: $\bar{\psi} = \phi^{-1}$
 (1)6. LET: $\bar{\psi} : P \rightarrow A$ be the function $\bar{\psi}(p) = \sup\{\psi(q) \mid q \in Q, q < p\}$
 (1)7. $\bar{\psi}$ is monotone and extends ψ
 PROOF: Similar.
 (1)8. $\bar{\psi} \circ \bar{\phi} : A \rightarrow A$ is monotone and the identity on B .
 (1)9. $\bar{\psi} \circ \bar{\phi} = \text{id}_A$
 PROOF: Proposition 5.1.20.
 (1)10. $\bar{\phi} \circ \bar{\psi} = \text{id}_B$
 PROOF: Proposition 5.1.20.
 (1)11. If $\phi^* : A \cong P$ is any order isomorphism that extends ϕ then $\phi^* = \bar{\phi}$.
 (2)1. LET: $a \in A$
 PROVE: $\phi^*(a) = \sup \phi(S(a))$
 (2)2. $\phi^*(a)$ is an upper bound for $\phi(S(a))$
 (2)3. LET: u be any upper bound for $\phi(S(a))$
 PROVE: $\phi^*(a) \leq u$
 (2)4. ASSUME: for a contradiction $u < \phi^*(a)$
 (2)5. PICK $q \in Q$ such that $u < q < \phi^*(a)$
 (2)6. PICK $b \in B$ such that $q = \phi(b)$
 (2)7. $b < a$
 (2)8. $b \in S(a)$
 (2)9. $q = \phi(b) \leq u$
 (2)10. Q.E.D.
 PROOF: This is a contradiction.

□

Definition 5.1.22 (Initial Segment). Let \leq be a partial order on \mathbf{A} and $t \in A$. The *initial segment* up to t is the class

$$\text{seg } t := \{x \in \mathbf{A} \mid x < t\} .$$

Definition 5.1.23 (Lexicographic Ordering). Let \mathbf{R} be a partial order on \mathbf{A} and \mathbf{S} a partial order on \mathbf{B} . The *lexicographic ordering* \leq on $\mathbf{A} \times \mathbf{B}$ is defined by:

$$(a, b) \leq (a', b') \Leftrightarrow (a\mathbf{R}a' \wedge a \neq a') \vee (a = a' \wedge b\mathbf{S}b') .$$

Proposition Schema 5.1.24 (Z). *For any classes \mathbf{A} , \mathbf{B} , \mathbf{R} and \mathbf{S} , the following is a theorem:*

If \mathbf{R} is a partial order on \mathbf{A} and \mathbf{S} is a partial order on \mathbf{B} then the lexicographic ordering on $\mathbf{A} \times \mathbf{B}$ is a partial order.

PROOF:

$\langle 1 \rangle 1$. LET: \leq be the lexicographic ordering on $\mathbf{A} \times \mathbf{B}$

$\langle 1 \rangle 2$. \leq is reflexive.

PROOF: For any $a \in \mathbf{A}$ and $b \in \mathbf{B}$ we have $a = a$ and $b\mathbf{S}b$, so $(a, b) \leq (a, b)$.

$\langle 1 \rangle 3$. \leq is antisymmetric.

$\langle 2 \rangle 1$. ASSUME: $(a, b) \leq (a', b')$ and $(a', b') \leq (a, b)$

$\langle 2 \rangle 2$. $(a\mathbf{R}a' \wedge a \neq a') \vee (a = a' \wedge b\mathbf{S}b')$

$\langle 2 \rangle 3$. $(a'\mathbf{R}a \wedge a' \neq a) \vee (a' = a \wedge b\mathbf{S}b')$

$\langle 2 \rangle 4$. CASE: $a = a'$

PROOF: Then $b\mathbf{S}b'$ and $b'\mathbf{S}b$ hence $b = b'$ and $(a, b) = (a', b')$.

$\langle 2 \rangle 5$. CASE: $a \neq a'$

PROOF: Then $a\mathbf{R}a'$ and $a'\mathbf{R}a$ hence $a = a'$ which is a contradiction.

$\langle 1 \rangle 4$. \leq is transitive.

$\langle 2 \rangle 1$. ASSUME: $(a_1, b_1) \leq (a_2, b_2) \leq (a_3, b_3)$

$\langle 2 \rangle 2$. $(a_1\mathbf{R}a_2 \wedge a_1 \neq a_2) \vee (a_1 = a_2 \wedge b_1\mathbf{S}b_2)$

$\langle 2 \rangle 3$. $(a_2\mathbf{R}a_3 \wedge a_2 \neq a_3) \vee (a_2 = a_3 \wedge b_2\mathbf{S}b_3)$

$\langle 2 \rangle 4$. CASE: $a_1\mathbf{R}a_2, a_1 \neq a_2, a_2\mathbf{R}a_3, a_2 \neq a_3$

$\langle 3 \rangle 1$. $a_1\mathbf{R}a_3$

PROOF: Since \mathbf{R} is transitive.

$\langle 3 \rangle 2$. $a_1 \neq a_3$

PROOF: If $a_1 = a_3$ then $a_1\mathbf{R}a_2$ and $a_2\mathbf{R}a_1$ so $a_1 = a_2$ which is a contradiction.

$\langle 2 \rangle 5$. CASE: $a_1\mathbf{R}a_2, a_1 \neq a_2, a_2 = a_3, b_2\mathbf{S}b_3$

PROOF: Then $a_1\mathbf{R}a_3$ and $a_1 \neq a_3$.

$\langle 2 \rangle 6$. CASE: $a_1 = a_2, b_1\mathbf{S}b_2, a_2\mathbf{R}a_3, a_2 \neq a_3$

PROOF: Then $a_1\mathbf{R}a_3$ and $a_1 \neq a_3$.

$\langle 2 \rangle 7$. CASE: $a_1 = a_2, b_1\mathbf{S}b_2, a_2 = a_3, b_2\mathbf{S}b_3$

PROOF: Then $a_1 = a_3$ and $b_1\mathbf{S}b_3$.

□

5.2 Linear Orders

Definition 5.2.1 (Linear Ordering). Let \mathbf{A} be a class. A *linear ordering* or *total ordering* on \mathbf{A} is a partial ordering \leq on \mathbf{A} that is *total*, i.e.

$$\forall x, y \in \mathbf{A}. x \leq y \vee y \leq x$$

We often use the symbol $<$ for a linear ordering, and then write $x < y$ for $(x, y) \in <$.

Proposition Schema 5.2.2 (Trichotomy (Z)). *For any classes \mathbf{A} and \leq , the following is a theorem:*

Assume \leq be a linear ordering on \mathbf{A} . For any $x, y \in \mathbf{A}$, exactly one of $x < y$, $x = y$, $y < x$ holds.

PROOF: Immediate from definitions. \square

Proposition Schema 5.2.3 (Z). *For any classes \mathbf{A} and $<$, the following is a theorem:*

Let $<$ be a transitive relation on \mathbf{A} that satisfies trichotomy. Define \leq on \mathbf{A} by $x \leq y$ iff $x < y$ or $x = y$. Then \leq is a linear ordering on \mathbf{A} and $x < y$ iff $x \leq y$ and $x \neq y$.

PROOF:

$\langle 1 \rangle 1.$ \leq is reflexive.

PROOF: By definition we have $\forall x \in \mathbf{A}. x \leq x$.

$\langle 1 \rangle 2.$ \leq is antisymmetric.

$\langle 2 \rangle 1.$ ASSUME: $x \leq y$ and $y \leq x$

$\langle 2 \rangle 2.$ $x < y$ or $x = y$

$\langle 2 \rangle 3.$ $y < x$ or $y = x$

$\langle 2 \rangle 4.$ We cannot have $x < y$ and $y < x$

PROOF: Trichotomy.

$\langle 2 \rangle 5.$ $x = y$

$\langle 1 \rangle 3.$ \leq is transitive.

$\langle 2 \rangle 1.$ ASSUME: $x \leq y$ and $y \leq z$

$\langle 2 \rangle 2.$ $x < y$ or $x = y$

$\langle 2 \rangle 3.$ $y < z$ or $y = z$

$\langle 2 \rangle 4.$ CASE: $x < y$ and $y < z$

PROOF: Then $x < z$ by transitivity, so $x \leq z$.

$\langle 2 \rangle 5.$ CASE: $x = y$

PROOF: Then we have $y \leq z$ and so $x \leq z$.

$\langle 2 \rangle 6.$ CASE: $y = z$

PROOF: Then we have $x \leq y$ and so $x \leq z$.

$\langle 1 \rangle 4.$ \leq is total.

PROOF: Immediate from trichotomy.

\square

Proposition Schema 5.2.4 (Z). *For any classes \mathbf{A} and \mathbf{R} , the following is a theorem:*

If \mathbf{R} is a linear ordering on \mathbf{A} then \mathbf{R}^{-1} is also a linear ordering on \mathbf{A} .

PROOF:

$\langle 1 \rangle 1.$ \mathbf{R}^{-1} is a partial order on \mathbf{A} .

PROOF: Proposition 5.1.2.

$\langle 1 \rangle 2.$ \mathbf{R}^{-1} is total.

- $\langle 2 \rangle 1.$ LET: $x, y \in \mathbf{A}$
- $\langle 2 \rangle 2.$ $x\mathbf{R}y$ or $y\mathbf{R}x$.
- $\langle 2 \rangle 3.$ $y\mathbf{R}^{-1}x$ or $x\mathbf{R}^{-1}y$.

□

Proposition Schema 5.2.5 (Z). *For any classes $\mathbf{A}, \mathbf{B}, \mathbf{F}, \mathbf{R}, \mathbf{S}$, the following is a theorem:*

Assume \mathbf{R} is a linear order on \mathbf{A} , \mathbf{S} is a partial order on \mathbf{B} , and $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$. If \mathbf{F} is strictly monotone then it is injective.

PROOF:

- $\langle 1 \rangle 1.$ LET: $x, y \in \mathbf{A}$
- $\langle 1 \rangle 2.$ ASSUME: $x \neq y$
PROVE: $\mathbf{F}(x) \neq \mathbf{F}(y)$
- $\langle 1 \rangle 3.$ ASSUME: w.l.o.g. $x\mathbf{R}y$
PROOF: \mathbf{R} is total.
- $\langle 1 \rangle 4.$ $\mathbf{F}(x)\mathbf{S}\mathbf{F}(y)$ and $\mathbf{F}(x) \neq \mathbf{F}(y)$
PROOF: \mathbf{F} is strictly monotone.

□

Proposition Schema 5.2.6 (Z). *For any classes \mathbf{A}, \mathbf{B} , \leq , \preceq and \mathbf{F} , the following is a theorem:*

Assume \leq is a linear order on \mathbf{A} and \preceq is a linear order on \mathbf{B} . Assume $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and \mathbf{F} is strictly monotone. For all $x, y \in \mathbf{A}$, if $\mathbf{F}(x) \prec \mathbf{F}(y)$ then $x < y$.

PROOF:

- $\langle 1 \rangle 1.$ $\mathbf{F}(x) \neq \mathbf{F}(y)$ and $\mathbf{F}(y) \not\prec \mathbf{F}(x)$
PROOF: Trichotomy.
- $\langle 1 \rangle 2.$ $x \neq y$ and $y \not\prec x$
PROOF: \mathbf{F} is strictly monotone.
- $\langle 1 \rangle 3.$ $x < y$
PROOF: Trichotomy.

□

Corollary Schema 5.2.6.1 (Z). *For any classes \mathbf{A}, \mathbf{B} , \leq , \preceq and \mathbf{F} , the following is a theorem:*

Assume \leq is a linear order on \mathbf{A} and \preceq is a linear order on \mathbf{B} . Assume $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ and \mathbf{F} is strictly monotone. Then \mathbf{F} is an order isomorphism.

Proposition Schema 5.2.7 (Z). *For any classes $\mathbf{A}, \mathbf{B}, \mathbf{F}$ and \mathbf{S} , the following is a theorem:*

Assume \mathbf{S} is a linear order on \mathbf{B} and $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$. Define \mathbf{R} on \mathbf{A} by $x\mathbf{R}y$ if and only if $\mathbf{F}(x)\mathbf{S}\mathbf{F}(y)$. Then \mathbf{R} is a linear order on \mathbf{A} .

PROOF:

- $\langle 1 \rangle 1.$ \mathbf{R} is a partial order on \mathbf{A} .
PROOF: Proposition 5.1.3.

$\langle 1 \rangle 2$. \mathbf{R} is total.

PROOF: For all $x, y \in \mathbf{A}$ we have $\mathbf{F}(x)\mathbf{SF}(y)$ or $\mathbf{F}(y)\mathbf{SF}(x)$.

□

Corollary Schema 5.2.7.1 (Z). *For any classes \mathbf{A} , \mathbf{B} and \mathbf{R} , the following is a theorem:*

Assume \mathbf{R} be a linear order on \mathbf{A} and $\mathbf{B} \subseteq \mathbf{A}$. Then $\mathbf{R} \cap \mathbf{B}^2$ is a linear order on \mathbf{B} .

Proposition Schema 5.2.8 (Z). *For any classes \mathbf{A} , \mathbf{B} , \mathbf{R} and \mathbf{S} , the following is a theorem:*

Assume \mathbf{R} is a linear order on \mathbf{A} and \mathbf{S} is a linear order on \mathbf{B} . Then the lexicographic ordering is a linear order on $\mathbf{A} \times \mathbf{B}$.

PROOF:

$\langle 1 \rangle 1$. LET: \leq be the lexicographic order on $\mathbf{A} \times \mathbf{B}$

$\langle 1 \rangle 2$. \leq is a partial order.

PROOF: Proposition 5.1.24.

$\langle 1 \rangle 3$. \leq is total.

$\langle 2 \rangle 1$. LET: $a, a' \in \mathbf{A}$ and $b, b' \in \mathbf{B}$

$\langle 2 \rangle 2$. CASE: $a\mathbf{R}a'$ and $a \neq a'$

PROOF: Then $(a, b) \leq (a', b')$.

$\langle 2 \rangle 3$. CASE: $a = a'$

PROOF: We have $b\mathbf{S}b'$ or $b'\mathbf{S}b$, so $(a, b) \leq (a', b')$ or $(a', b') \leq (a, b)$.

$\langle 2 \rangle 4$. CASE: $a'\mathbf{R}a$ and $a \neq a'$

PROOF: Then $(a', b') \leq (a, b)$.

□

5.3 Well Orderings

Definition 5.3.1 (Well Ordering). A *well ordering* on a class \mathbf{A} is a well-founded linear ordering on \mathbf{A} .

Proposition 5.3.2 (Z). *Let S be a well ordering of the set B and $f : A \rightarrow B$ a function. Define R on A by xRy if and only if $F(x)SF(y)$. Then R well orders A .*

PROOF:

$\langle 1 \rangle 1$. R linearly orders A .

PROOF: Proposition 5.2.7.

$\langle 1 \rangle 2$. Every nonempty subset of A has a least element.

$\langle 2 \rangle 1$. LET: C be a nonempty subset of A .

$\langle 2 \rangle 2$. LET: y be the least element of $f(C)$.

$\langle 2 \rangle 3$. PICK $x \in C$ such that $f(x) = y$.

$\langle 2 \rangle 4$. x is least in C .

□

Proposition Schema 5.3.3 (Z). *For any classes \mathbf{A} , \mathbf{B} and \mathbf{R} , the following is a theorem:*

Assume \mathbf{R} well orders \mathbf{B} and $\mathbf{A} \subseteq \mathbf{B}$. Then $\mathbf{R} \cap \mathbf{A}^2$ well orders \mathbf{A} .

PROOF:

$\langle 1 \rangle 1$. LET: $\mathbf{R}' = \mathbf{R} \cap \mathbf{A}^2$

$\langle 1 \rangle 2$. \mathbf{R}' linearly orders \mathbf{A} .

PROOF: Corollary 5.2.7.1.

$\langle 1 \rangle 3$. \mathbf{R}' is well founded.

PROOF: Proposition 2.4.4.

□

Proposition Schema 5.3.4 (ZFC). *For any classes \mathbf{A} , \mathbf{B} , \mathbf{F} and \mathbf{S} , the following is a theorem:*

Assume \mathbf{S} well orders \mathbf{B} and $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$. Define \mathbf{R} on \mathbf{A} by $x\mathbf{R}y$ if and only if $\mathbf{F}(x)\mathbf{S}\mathbf{F}(y)$. Then \mathbf{R} well orders \mathbf{A} .

PROOF:

$\langle 1 \rangle 1$. \mathbf{R} linearly orders \mathbf{A} .

PROOF: Proposition 5.2.7.

$\langle 1 \rangle 2$. For all $t \in \mathbf{A}$ we have $\{x \in \mathbf{A} \mid x\mathbf{R}t \wedge x \neq t\}$ is a set.

$\langle 2 \rangle 1$. LET: $t \in \mathbf{A}$

$\langle 2 \rangle 2$. LET: $S = \{y \in \mathbf{B} \mid y\mathbf{S}\mathbf{F}(t) \wedge y \neq \mathbf{F}(t)\}$

$\langle 2 \rangle 3$. LET: $P(x, y)$ be the property $\mathbf{F}(y) = x$

$\langle 2 \rangle 4$. For all $x \in S$ there exists at most one y such that $P(x, y)$

PROOF: \mathbf{F} is injective.

$\langle 2 \rangle 5$. LET: $T = \{y \mid \exists x \in S. P(x, y)\}$

PROOF: Axiom of Replacement.

$\langle 2 \rangle 6$. $T = \{x \in \mathbf{A} \mid x\mathbf{R}t \wedge x \neq t\}$

$\langle 1 \rangle 3$. Every nonempty subset of \mathbf{A} has a least element.

$\langle 2 \rangle 1$. LET: S be a nonempty subset of \mathbf{A} .

$\langle 2 \rangle 2$. $\mathbf{F}(S)$ is a nonempty subset of \mathbf{B}

PROOF: Axiom of Replacement.

$\langle 2 \rangle 3$. LET: y be the least element of $\mathbf{F}(S)$.

$\langle 2 \rangle 4$. PICK $x \in S$ such that $\mathbf{F}(x) = y$.

$\langle 2 \rangle 5$. x is least in S .

□

Proposition 5.3.5 (Z). *For any well ordered sets A and B , the lexicographic order well orders $A \times B$.*

PROOF:

$\langle 1 \rangle 1$. $A \times B$ is linearly ordered.

PROOF: Proposition 5.2.8.

$\langle 1 \rangle 2$. Every nonempty subset of $A \times B$ has a least element.

$\langle 2 \rangle 1$. LET: S be a nonempty subset of $A \times B$.

$\langle 2 \rangle 2$. LET: a be the least element of $\{x \in A \mid \exists y \in B. (x, y) \in S\}$.

$\langle 2 \rangle 3$. LET: b be the least element of $\{y \in B \mid (a, y) \in S\}$.

⟨2⟩4. (a, b) is least in S .
 \square

Definition 5.3.6 (End Extension). Let A and B be well ordered sets. Then B is an *end extension* of A iff $A \subseteq B$ and:

- Whenever $x, y \in A$ then $x \leq_A y$ iff $x \leq_B y$.
- Whenever $x \in A$ and $y \in B - A$ then $x < y$.

Theorem 5.3.7 (Z). Let \leq be a linear ordering on A . Assume that, for any $B \subseteq A$ such that $\forall t \in A. \text{seg } t \subseteq B \Rightarrow t \in B$, we have $B = A$. Then \leq is a well ordering on A .

PROOF:

- ⟨1⟩1. LET: $C \subseteq A$ be nonempty.
 ⟨1⟩2. LET: $B = \{t \in A \mid \forall x \in C. t < x\}$
 ⟨1⟩3. $B \cap C = \emptyset$
 ⟨1⟩4. $B \neq A$
 ⟨1⟩5. PICK $t \in A$ such that $\text{seg } t \subseteq B$ and $t \notin B$
 ⟨1⟩6. t is least in C .
 \square

Proposition Schema 5.3.8 (Z). For any classes $\mathbf{A}, \mathbf{B}, \mathbf{F}, \mathbf{G}, \leq$ and \preceq , the following is a theorem:

Assume \leq well orders \mathbf{A} and \preceq well orders \mathbf{B} . Assume \mathbf{F} and \mathbf{G} are order isomorphisms between \mathbf{A} and \mathbf{B} . Then $\mathbf{F} = \mathbf{G}$.

PROOF:

- ⟨1⟩1. For all $x \in \mathbf{A}$, if $\forall t < x. \mathbf{F}(t) = \mathbf{G}(t)$, then $\mathbf{F}(x) = \mathbf{G}(x)$
 ⟨2⟩1. LET: $x \in \mathbf{A}$
 ⟨2⟩2. ASSUME: $\forall t < x. \mathbf{F}(t) = \mathbf{G}(t)$
 ⟨2⟩3. $\mathbf{F}(\text{seg } x) = \mathbf{G}(\text{seg } x)$
 ⟨2⟩4. $\mathbf{F}(x)$ is the least element of $\mathbf{B} - \mathbf{F}(\text{seg } x)$
 ⟨2⟩5. $\mathbf{G}(x)$ is the least element of $\mathbf{B} - \mathbf{G}(\text{seg } x)$
 ⟨2⟩6. $\mathbf{F}(x) = \mathbf{G}(x)$
 ⟨1⟩2. $\forall x \in \mathbf{A}. \mathbf{F}(x) = \mathbf{G}(x)$

PROOF: Transfinite induction.

\square

Theorem 5.3.9 (ZFC). Let A and B be well ordered sets. Then one of the following holds: $A \cong B$; there exists $b \in B$ such that $A \cong \text{seg } b$; there exists $a \in A$ such that $\text{seg } a \cong B$.

PROOF:

- ⟨1⟩1. PICK e that is not in A or B .
 ⟨1⟩2. LET: $F : A \rightarrow B \cup \{e\}$ be the function defined by transfinite recursion thus:

$$F(t) = \begin{cases} \text{the least element of } B - F(\text{seg } t) & \text{if } B - F(\text{seg } t) \neq \emptyset \\ e & \text{if } B - F(\text{seg } t) = \emptyset \end{cases}$$

$\langle 2 \rangle 1$. LET: t be least such that $F(t) = e$

$\langle 2 \rangle 1$. LET: t be least such that $F(t) = e$

$$\langle 2 \rangle 2. \quad F \upharpoonright \text{seg } t : \text{seg } t \cong B$$

PROOF: We have $F : A \cong B$

PROOF: We have $F : A \cong B$

$\langle 2 \rangle 1$. LET: b be the least element of $B - \text{ran } F$

$\langle 2 \rangle 1$. LET: b be the least element of $B - \text{ran } F$

$\langle 2 \rangle 2. F : A \cong \text{seg } b$

1

Chapter 6

Ordinal Numbers

6.1 Ordinals

Definition 6.1.1 (Ordinal Number). An *ordinal (number)* is a transitive set α that is *well-ordered by* \in ; that is, such that $\{(x, y) \in \alpha^2 \mid x \in y \vee x = y\}$ well orders α .

Given $x, y \in \alpha$, we write $x < y$ iff $x \in y$, and $x \leq y$ iff $x \in y$ or $x = y$.

Let \mathbf{On} be the class of ordinal numbers. For $\alpha, \beta \in \mathbf{On}$, we write $\alpha < \beta$ iff $\alpha \in \beta$, and $\alpha \leq \beta$ iff $\alpha < \beta$ or $\alpha = \beta$.

Proposition 6.1.2 (Z). *For any ordinal numbers α and β , if $\alpha \cong \beta$ then $\alpha = \beta$.*

PROOF:

$\langle 1 \rangle 1$. LET: $f : \alpha \cong \beta$

$\langle 1 \rangle 2$. For all $x \in \alpha$, if $\forall t < x. f(t) = t$ then $f(x) = x$

$\langle 2 \rangle 1$. $f(x) \subseteq x$

$\langle 3 \rangle 1$. LET: $y \in f(x)$

$\langle 3 \rangle 2$. $y \in \beta$

$\langle 3 \rangle 3$. PICK $t \in \alpha$ such that $f(t) = y$

PROOF: f is surjective.

$\langle 3 \rangle 4$. $f(t) \in f(x)$

$\langle 3 \rangle 5$. $t \in x$

PROOF: Since f is an order isomorphism.

$\langle 3 \rangle 6$. $f(t) = t$

PROOF: Induction hypothesis.

$\langle 3 \rangle 7$. $y = t$

$\langle 3 \rangle 8$. $y \in x$

$\langle 2 \rangle 2$. $x \subseteq f(x)$

$\langle 3 \rangle 1$. LET: $t \in x$

$\langle 3 \rangle 2$. $f(t) \in f(x)$

$\langle 3 \rangle 3$. $f(t) = t$

$\langle 3 \rangle 4$. $t \in f(x)$

$\langle 1 \rangle 3. \forall x \in \alpha. f(x) = x$

PROOF: Transfinite induction.

$\langle 1 \rangle 4. \alpha = \beta$

PROOF: Since $\beta = \{f(t) \mid t \in \alpha\} = \{t \mid t \in \alpha\} = \alpha$.

□

Theorem 6.1.3 (ZFC). *Every well-ordered set is isomorphic to a unique ordinal.*

PROOF:

$\langle 1 \rangle 1.$ For any well-ordered set A , there exists an ordinal α such that $A \cong \alpha$.

$\langle 2 \rangle 1.$ LET: A be a well-ordered set.

$\langle 2 \rangle 2.$ Define the function E on A by transfinite recursion thus:

$$E(t) = \{E(x) \mid x < t\} \quad (t \in A) .$$

$\langle 2 \rangle 3.$ LET: $\alpha = \{E(x) \mid x \in A\}$

$\langle 2 \rangle 4.$ α is an ordinal.

$\langle 3 \rangle 1.$ α is a transitive set.

$\langle 4 \rangle 1.$ LET: $x \in y \in \alpha$

$\langle 4 \rangle 2.$ PICK $t \in A$ such that $y = E(t)$

$\langle 4 \rangle 3.$ $x \in E(t) = \{E(s) \mid s < t\}$

$\langle 4 \rangle 4.$ PICK $s < t$ such that $x = E(s)$

$\langle 4 \rangle 5.$ $x \in \alpha$

$\langle 3 \rangle 2.$ α is well-ordered by \in .

$\langle 4 \rangle 1.$ LET: $< = \{(x, y) \in \alpha \mid x \in y\}$

$\langle 4 \rangle 2.$ $<$ is transitive.

$\langle 5 \rangle 1.$ LET: $x, y, z \in \alpha$ with $x \in y \in z$

$\langle 5 \rangle 2.$ PICK $t \in A$ such that $z = E(t)$

$\langle 5 \rangle 3.$ PICK $s \in A$ such that $s < t$ and $y = E(s)$

$\langle 5 \rangle 4.$ PICK $r \in A$ such that $r < s$ and $x = E(r)$

$\langle 5 \rangle 5.$ $r < t$

$\langle 5 \rangle 6.$ $x \in z$

$\langle 4 \rangle 3.$ $<$ satisfies trichotomy.

$\langle 5 \rangle 1.$ LET: $x, y \in \alpha$

$\langle 5 \rangle 2.$ PICK $s, t \in A$ such that $E(s) = x$ and $E(t) = y$

$\langle 5 \rangle 3.$ Exactly one of $s < t$, $s = t$, $t < s$ holds.

$\langle 5 \rangle 4.$ CASE: $s < t$

$\langle 6 \rangle 1.$ $x \in y$

$\langle 6 \rangle 2.$ $x \neq y$ and $y \notin x$

PROOF: Axiom of Regularity.

$\langle 5 \rangle 5.$ CASE: $s = t$

$\langle 6 \rangle 1.$ $x = y$

$\langle 6 \rangle 2.$ $x \notin y$ and $y \notin x$

PROOF: Axiom of Regularity.

$\langle 5 \rangle 6.$ CASE: $t < s$

PROOF: Similar to $\langle 5 \rangle 4$.

$\langle 4 \rangle 4.$ \leq is a linear order on α .

PROOF: Proposition 5.2.3.

- ⟨4⟩5. Every nonempty subset of α has a least element.
 - ⟨5⟩1. LET: S be a nonempty subset of α
 - ⟨5⟩2. LET: $T = \{x \in A \mid E(x) \in S\}$
 - ⟨5⟩3. LET: t be the least element of T .
PROVE: $E(t)$ is least in S
 - ⟨5⟩4. LET: $y \in S$
 - ⟨5⟩5. PICK $s \in T$ such that $E(s) = y$
 - ⟨5⟩6. $t \leq s$
 - ⟨5⟩7. $x \leq y$
- ⟨2⟩5. E is surjective.
PROOF: By definition of α .
- ⟨2⟩6. E is strictly monotone.
PROOF: If $s < t$ then $E(s) \in E(t)$ by definition of $E(t)$.
- ⟨2⟩7. Q.E.D.
PROOF: Corollary 5.2.6.1.
- ⟨1⟩2. For any ordinals α and β , if $\alpha \cong \beta$ then $\alpha = \beta$.
PROOF: Proposition 6.1.2.

□

Proposition 6.1.4 (Z). *The class **On** is a transitive class. That is, every element of an ordinal is an ordinal.*

PROOF:

- ⟨1⟩1. LET: α be an ordinal.
- ⟨1⟩2. LET: $\beta \in \alpha$
- ⟨1⟩3. β is a transitive set.
 - ⟨2⟩1. LET: $x \in y \in \beta$
 - ⟨2⟩2. $y \in \alpha$
PROOF: α is transitive.
 - ⟨2⟩3. $x \in \alpha$
PROOF: α is transitive.
 - ⟨2⟩4. $x \in \beta$
PROOF: Since $\{(x, y) \in \alpha^2 \mid x \in y\}$ is transitive.
- ⟨1⟩4. β is well ordered by \in .
PROOF: By Proposition 5.3.3.

□

Proposition 6.1.5 (ZFC). *Given two ordinal numbers α, β , exactly one of $\alpha \in \beta$, $\alpha = \beta$, $\beta \in \alpha$ holds.*

PROOF:

- ⟨1⟩1. At most one holds.
PROOF: Since every ordinal is a transitive set and we never have $\alpha \in \alpha$.
- ⟨1⟩2. At least one holds.
 - ⟨2⟩1. Either $\alpha \cong \beta$ or $\exists t \in \beta. \alpha \cong \text{seg } t$ or $\exists t \in \alpha. \text{seg } t \cong \beta$.
 - ⟨2⟩2. CASE: $\alpha \cong \beta$
PROOF: Then $\alpha = \beta$ by Proposition 6.1.2.

$\langle 2 \rangle 3$. CASE: There exists $t \in \beta$ such that $\alpha \cong \text{seg } t$

$\langle 3 \rangle 1$. t is an ordinal number.

PROOF: Proposition 6.1.4.

$\langle 3 \rangle 2$. $t = \text{seg } t$

$\langle 4 \rangle 1$. $t \subseteq \text{seg } t$

$\langle 5 \rangle 1$. LET: $s \in t$

$\langle 5 \rangle 2$. $s \in \beta$

PROOF: β is a transitive set.

$\langle 5 \rangle 3$. $s \in \text{seg } t$

$\langle 4 \rangle 2$. $\text{seg } t \subseteq t$

PROOF: Immediate from definitions.

$\langle 3 \rangle 3$. $\alpha = t$

PROOF: Proposition 6.1.2.

$\langle 3 \rangle 4$. $\alpha \in \beta$

$\langle 2 \rangle 4$. CASE: There exists $t \in \alpha$ such that $\text{seg } t \cong \beta$

PROOF: $\beta \in \alpha$ similarly.

□

Proposition 6.1.6 (Z). *Any nonempty set S of ordinal numbers has a least element.*

PROOF:

$\langle 1 \rangle 1$. PICK $\beta \in S$

$\langle 1 \rangle 2$. CASE: $\beta \cap S = \emptyset$

PROOF: Then β is least in S .

$\langle 1 \rangle 3$. CASE: $\beta \cap S \neq \emptyset$

PROOF: The least element of $\beta \cap S$ is least in S .

□

Theorem 6.1.7 (ZFC). *The class \mathbf{On} is well ordered by \in .*

PROOF:

$\langle 1 \rangle 1$. LET: $\mathbf{E} = \{(x, y) \in \mathbf{On}^2 \mid x \in y\}$

$\langle 1 \rangle 2$. \mathbf{E} is transitive.

PROOF: If $\alpha \in \beta \in \gamma$ then $\alpha \in \gamma$ because every ordinal is a transitive set.

$\langle 1 \rangle 3$. \mathbf{E} satisfies trichotomy.

PROOF: Proposition 6.1.5.

$\langle 1 \rangle 4$. \mathbf{E} linearly orders \mathbf{On} .

PROOF: Proposition 5.2.3.

$\langle 1 \rangle 5$. \mathbf{E} is well founded.

PROOF: Proposition 2.4.2.

□

Corollary 6.1.7.1 (Burali-Forti Paradox (ZFC)). *The class \mathbf{On} is a proper class.*

PROOF: If it were a set, it would be a transitive set well-ordered by \in , and hence a member of itself, contradicting Proposition 1.5.3.

Proposition 6.1.8 (ZFC). *Any transitive set of ordinal numbers is an ordinal number.*

PROOF: It is well-ordered by \in by Proposition 5.3.3 and Theorem 6.1.7. \square

Proposition 6.1.9 (Z). *\emptyset is an ordinal number.*

PROOF: Vacuously, it is a transitive set well-ordered by \in .

Definition 6.1.10. We define $0 = \emptyset$.

Proposition 6.1.11 (ZFC). *If A is a set of ordinal numbers then $\bigcup A$ is an ordinal number.*

PROOF:

$\langle 1 \rangle 1.$ $\bigcup A$ is a transitive set.

PROOF: Proposition 1.6.3.

$\langle 1 \rangle 2.$ $\bigcup A$ is a set of ordinals.

PROOF: Proposition 6.1.4.

$\langle 1 \rangle 3.$ Q.E.D.

PROOF: Proposition 6.1.8.

\square

Corollary 6.1.11.1 (ZFC). *The poset \mathbf{On} is complete.*

PROOF: For any nonempty set A of ordinals, $\bigcup A$ is its supremum. \square

Proposition 6.1.12 (ZFC). *Let α be an ordinal and $S \subseteq \alpha$. Then S is well-ordered by \in and the ordinal of (S, \in) is $\leq \alpha$.*

PROOF:

$\langle 1 \rangle 1.$ S is well ordered by \in .

$\langle 1 \rangle 2.$ LET: β be the ordinal of (S, \in)

$\langle 1 \rangle 3.$ LET: $E : S \approx \beta$ be the unique isomorphism.

$\langle 1 \rangle 4.$ $\forall \gamma \in S. E(\gamma) \leq \gamma$

$\langle 2 \rangle 1.$ LET: $\gamma \in S$

$\langle 2 \rangle 2.$ ASSUME: as transfinite induction hypothesis $\forall \delta < \gamma. E(\delta) \leq \delta$

$\langle 2 \rangle 3.$ $E(\gamma)$ is the least element of β that is greater than $E(\delta)$ for all $\delta < \gamma$

$\langle 2 \rangle 4.$ γ is greater than $E(\delta)$ for all $\delta < \gamma$

$\langle 2 \rangle 5.$ $E(\gamma) \leq \gamma$

$\langle 1 \rangle 5.$ $\beta \leq \alpha$

$\langle 2 \rangle 1.$ $\forall \gamma < \beta. \gamma < \alpha$

$\langle 3 \rangle 1.$ LET: $\gamma < \beta$

$\langle 3 \rangle 2.$ PICK $\delta \in S$ such that $E(\delta) = \gamma$

$\langle 3 \rangle 3.$ $\gamma = E(\delta) \leq \delta < \alpha$

\square

Proposition 6.1.13 (ZFC). *Let α be a set. Then the following are equivalent.*

1. α is an ordinal.

2. α is a transitive set and, for all $x, y \in \alpha$, either $x = y$ or $x \in y$ or $y \in x$.

3. α is a transitive set of transitive sets.

PROOF:

$\langle 1 \rangle 1. 1 \Rightarrow 2$

PROOF: Immediate from definitions.

$\langle 1 \rangle 2. 2 \Rightarrow 3$

$\langle 2 \rangle 1.$ ASSUME: α is a transitive set and, for all $x, y \in \alpha$, either $x = y$ or $x \in y$ or $y \in x$

$\langle 2 \rangle 2.$ LET: $z \in \alpha$

PROVE: z is transitive.

$\langle 2 \rangle 3.$ LET: $x \in y \in z$

$\langle 2 \rangle 4.$ $y \in \alpha$

$\langle 2 \rangle 5.$ $x \in \alpha$

$\langle 2 \rangle 6.$ Either $x = z$ or $x \in z$ or $z \in x$

$\langle 2 \rangle 7.$ $x \neq z$

PROOF: We cannot have $x \in y \in x$ by the Axiom of Regularity.

$\langle 2 \rangle 8.$ $z \notin x$

PROOF: We cannot have $x \in y \in z \in x$ by the Axiom of Regularity.

$\langle 1 \rangle 3. 3 \Rightarrow 1$

$\langle 2 \rangle 1.$ LET: x be a transitive set of transitive sets.

$\langle 2 \rangle 2.$ ASSUME: as \in -induction hypothesis that, for all $y \in x$, if y is a transitive set of transitive sets then y is a transitive set of ordinals.

$\langle 2 \rangle 3.$ Every element of x is an ordinal.

$\langle 3 \rangle 1.$ LET: $y \in x$

$\langle 3 \rangle 2.$ y is transitive.

$\langle 3 \rangle 3.$ Every element of y is transitive.

PROOF: Since every element of y is an element of x , because x is transitive.

$\langle 3 \rangle 4.$ y is an ordinal.

PROOF: $\langle 2 \rangle 2$

$\langle 2 \rangle 4.$ Q.E.D.

PROOF: Proposition 6.1.8.

□

Lemma 6.1.14 (Z). *Let A and B be well-ordered sets. If B is an end extension of A then the ordinal of A is \leq the ordinal of B .*

PROOF:

$\langle 1 \rangle 1.$ LET: α be the ordinal of A and β the ordinal of B .

$\langle 1 \rangle 2.$ LET: $E_A : A \cong \alpha$ and $E_B : B \cong \beta$ be the canonical isomorphisms.

$\langle 1 \rangle 3.$ $\forall a \in A. E_A(a) = E_B(a)$

$\langle 2 \rangle 1.$ LET: $a \in A$

$\langle 2 \rangle 2.$ ASSUME: as transfinite induction hypothesis $\forall x < a. E_A(x) = E_B(x)$

$\langle 2 \rangle 3.$ $E_A(a)$ is the least ordinal that is greater than $E_A(x)$ for all $x < a$

$\langle 2 \rangle 4.$ $E_B(a)$ is the least ordinal that is greater than $E_B(x)$ for all $x < b$

- $\langle 2 \rangle 5. \{x \in A \mid x <_A a\} = \{x \in B \mid x <_B a\}$
- $\langle 2 \rangle 6. E_A(a) = E_B(a)$
- $\langle 1 \rangle 4. \alpha \subseteq \beta$
- $\langle 1 \rangle 5. \alpha \leq \beta$

□

Lemma 6.1.15. *Let \mathcal{C} be a set of well ordered sets such that, for any $A, B \in \mathcal{C}$, we have that one of A and B is an end extension of the other. Let $W = \bigcup \mathcal{C}$ under $x \leq y$ iff there exists $A \in \mathcal{C}$ such that $x, y \in A$ and $x \leq y$. Then W is a well ordered set whose ordinal is the supremum of the ordinals of the members of \mathcal{C} .*

PROOF:

- $\langle 1 \rangle 1. \leq$ is reflexive on W .
 - $\langle 2 \rangle 1.$ LET: $x \in W$
 - $\langle 2 \rangle 2.$ PICK $A \in \mathcal{C}$ such that $x \in A$.
 - $\langle 2 \rangle 3. x \leq x$
- $\langle 1 \rangle 2. \leq$ is antisymmetric on W .
 - $\langle 2 \rangle 1.$ LET: $x, y \in W$
 - $\langle 2 \rangle 2.$ ASSUME: $x \leq y$ and $y \leq x$
 - $\langle 2 \rangle 3.$ PICK $A \in \mathcal{C}$ such that $x, y \in A$ and $x \leq_A y$, and $B \in \mathcal{C}$ such that $x, y \in B$ and $y \leq_B x$
 - $\langle 2 \rangle 4.$ ASSUME: w.l.o.g. B is an end extension of A
 - $\langle 2 \rangle 5. x \leq_B y$ and $y \leq_B x$
 - $\langle 2 \rangle 6. x = y$
- $\langle 1 \rangle 3. \leq$ is transitive on W .
 - $\langle 2 \rangle 1.$ ASSUME: $x \leq y \leq z$
 - $\langle 2 \rangle 2.$ PICK $A, B \in \mathcal{C}$ such that $x \leq_A y$ and $y \leq_B z$
 - $\langle 2 \rangle 3.$ CASE: A is an end extension of B .
 - $\langle 3 \rangle 1. x \leq_A y$ and $y \leq_A z$
 - $\langle 3 \rangle 2. x \leq_A z$
 - $\langle 3 \rangle 3. x \leq z$
 - $\langle 2 \rangle 4.$ CASE: B is an end extension of A .

PROOF: Similar.
- $\langle 1 \rangle 4. \leq$ is total on W .
 - $\langle 2 \rangle 1.$ LET: $x, y \in W$
 - $\langle 2 \rangle 2.$ PICK $A, B \in \mathcal{C}$ such that $x \in A$ and $y \in B$
 - $\langle 2 \rangle 3.$ ASSUME: w.l.o.g. B is an end extension of A
 - $\langle 2 \rangle 4. x \leq_B y$ or $y \leq_B x$
 - $\langle 2 \rangle 5. x \leq_W y$ or $y \leq_W x$
- $\langle 1 \rangle 5.$ Every nonempty subset of W has a least element.
 - $\langle 2 \rangle 1.$ LET: S be a nonempty subset of W
 - $\langle 2 \rangle 2.$ PICK $s \in S$
 - $\langle 2 \rangle 3.$ PICK $A \in \mathcal{C}$ such that $s \in A$
 - $\langle 2 \rangle 4.$ LET: a be the \leq_A -least element of $S \cap A$

PROVE: a is least in S
 - $\langle 2 \rangle 5.$ LET: $x \in S$

- PROVE: $a \leq x$
- $\langle 2 \rangle 6$. PICK $B \in \mathcal{C}$ such that $x \in B$
 - $\langle 2 \rangle 7$. CASE: A is an end extension of B
 - $\langle 3 \rangle 1$. $a \leq_A x$
 - $\langle 3 \rangle 2$. $a \leq x$
 - $\langle 2 \rangle 8$. CASE: B is an end extension of A
 - $\langle 3 \rangle 1$. CASE: $x \in A$
 - $\langle 4 \rangle 1$. $a \leq_A x$
 - $\langle 4 \rangle 2$. $a \leq x$
 - $\langle 3 \rangle 2$. CASE: $x \in B - A$
 - $\langle 4 \rangle 1$. $a \leq_B x$
 - $\langle 4 \rangle 2$. $a \leq x$
 - $\langle 1 \rangle 6$. For all $A \in \mathcal{C}$, W is an end extension of A .
 - $\langle 2 \rangle 1$. For all $x, y \in A$, we have $x \leq_A y$ if and only if $x \leq_W y$
 - $\langle 3 \rangle 1$. LET: $x, y \in A$
 - $\langle 3 \rangle 2$. If $x \leq_A y$ then $x \leq_W y$
PROOF: Immediate from definitions.
 - $\langle 3 \rangle 3$. If $x \leq_W y$ then $x \leq_A y$
 - $\langle 4 \rangle 1$. ASSUME: $x \leq_W y$
 - $\langle 4 \rangle 2$. PICK $B \in \mathcal{C}$ such that $x \leq_B y$
 - $\langle 4 \rangle 3$. CASE: A is an end extension of B
PROOF: Then $x \leq_A y$.
 - $\langle 4 \rangle 4$. CASE: B is an end extension of A
PROOF: Then $x \leq_A y$.
 - $\langle 2 \rangle 2$. For all $x \in A$ and $y \in W - A$ we have $x < y$
 - $\langle 3 \rangle 1$. LET: $x \in A$ and $y \in W - A$
 - $\langle 3 \rangle 2$. PICK $B \in \mathcal{C}$ such that $y \in B$
 - $\langle 3 \rangle 3$. B is an end extension of A
 - $\langle 3 \rangle 4$. $x <_B y$
 - $\langle 3 \rangle 5$. $x <_W y$
 - $\langle 1 \rangle 7$. For all $A \in \mathcal{C}$, the ordinal of A is \leq the ordinal of W .
PROOF: Lemma 6.1.14.
 - $\langle 1 \rangle 8$. For any ordinal α , if for all $A \in \mathcal{C}$ the ordinal of A is $\leq \alpha$, then the ordinal of W is $\leq \alpha$.
 - $\langle 2 \rangle 1$. LET: α be an ordinal.
 - $\langle 2 \rangle 2$. ASSUME: for all $A \in \mathcal{C}$, the ordinal of A is $\leq \alpha$
 - $\langle 2 \rangle 3$. LET: β be the ordinal of W
 - $\langle 2 \rangle 4$. LET: $E : W \approx \beta$ be the canonical isomorphism.
 - $\langle 2 \rangle 5$. ASSUME: for a contradiction $\alpha < \beta$
 - $\langle 2 \rangle 6$. LET: $a \in W$ be the element with $E(a) = \alpha$
 - $\langle 2 \rangle 7$. PICK $A \in \mathcal{C}$ such that $a \in A$
 - $\langle 2 \rangle 8$. LET: γ be the ordinal of A and $E_A : A \cong \gamma$ be the canonical isomorphism.
phism.
 - $\langle 2 \rangle 9$. For all $x \in A$ we have $E_A(x) = E(x)$
PROOF: Transfinite induction on x .
 - $\langle 2 \rangle 10$. $E_A(a) = \alpha$

$\langle 2 \rangle 11.$ $\alpha < \gamma$

$\langle 2 \rangle 12.$ Q.E.D.

PROOF: This contradicts $\langle 2 \rangle 2$.

□

6.2 Successors

Definition 6.2.1 (Successor). The *successor* of a set a is the set $a^+ := a \cup \{a\}$.

Proposition 6.2.2 (Z). *A set a is a transitive set if and only if*

$$\bigcup(a^+) = a \text{ .}$$

PROOF:

$\langle 1 \rangle 1.$ If a is a transitive set then $\bigcup(a^+) = a$.

$\langle 2 \rangle 1.$ ASSUME: a is a transitive set.

$\langle 2 \rangle 2.$ $\bigcup(a^+) \subseteq a$

$\langle 3 \rangle 1.$ LET: $x \in \bigcup(a^+)$

PROVE: $x \in a$

$\langle 3 \rangle 2.$ PICK $y \in a^+$ such that $x \in y$.

$\langle 3 \rangle 3.$ $y \in a$ or $y = a$.

$\langle 3 \rangle 4.$ CASE: $y \in a$

PROOF: Then $x \in a$ because a is a transitive set.

$\langle 3 \rangle 5.$ CASE: $y = a$

PROOF: Then $x \in a$ immediately.

$\langle 2 \rangle 3.$ $a \subseteq \bigcup(a^+)$

PROOF: Since $a \in a^+$.

$\langle 1 \rangle 2.$ If $\bigcup(a^+) = a$ then a is a transitive set.

$\langle 2 \rangle 1.$ ASSUME: $\bigcup(a^+) = a$

$\langle 2 \rangle 2.$ $\bigcup a \subseteq a$

PROOF:

$$\begin{aligned} \bigcup a &\subseteq \bigcup(a^+) && \text{(Proposition 1.5.9)} \\ &= a && (\langle 2 \rangle 1) \end{aligned}$$

$\langle 2 \rangle 3.$ a is a transitive set.

PROOF: Proposition 1.6.2.

□

Proposition 6.2.3. *For any set a , we have a is a transitive set if and only if a^+ is a transitive set.*

PROOF:

$\langle 1 \rangle 1.$ If a is a transitive set then a^+ is a transitive set.

PROOF: If a is a transitive set then $\bigcup(a^+) = a \subseteq a^+$ by Proposition 6.2.2 and so a^+ is a transitive set.

$\langle 1 \rangle 2.$ If a^+ is a transitive set then a is a transitive set.

$\langle 2 \rangle 1.$ ASSUME: a^+ is a transitive set.

- $\langle 2 \rangle 2$. LET: $x \in y \in a$
- $\langle 2 \rangle 3$. $x \in y \in a^+$
- $\langle 2 \rangle 4$. $x \in a^+$
- PROOF: $\langle 2 \rangle 1$
- $\langle 2 \rangle 5$. $x \neq a$
- PROOF: From $\langle 2 \rangle 2$ and the Axiom of Regularity.
- $\langle 2 \rangle 6$. $x \in a$

□

Definition 6.2.4. We write 0 for \emptyset , 1 for \emptyset^+ , 2 for \emptyset^{++} , etc.

Proposition 6.2.5. For any set A we have $\mathcal{P}A \approx 2^A$.

PROOF: The function $H : \mathcal{P}A \rightarrow 2^A$ defined by $H(S)(a) = \{\emptyset\}$ if $a \in S$ and \emptyset if $a \notin S$ is a bijection. □

Proposition 6.2.6. For any ordinal number α we have α^+ is an ordinal number.

PROOF:

- $\langle 1 \rangle 1$. α^+ is a transitive set.
- PROOF: Proposition 6.2.3.
- $\langle 1 \rangle 2$. α^+ is well-ordered by \in .
- $\langle 2 \rangle 1$. For all $x, y, z \in \alpha^+$, if $x \in y \in z$ then $x \in z$
- $\langle 3 \rangle 1$. CASE: $z = \alpha$
- PROOF: Then $x \in \alpha$ since α is a transitive set.
- $\langle 3 \rangle 2$. CASE: $z \in \alpha$
- PROOF: Then $x \in z$ since α is well-ordered by \in .
- $\langle 2 \rangle 2$. For all $x, y \in \alpha^+$ we have $x \in y$ or $x = y$ or $y \in x$
- $\langle 3 \rangle 1$. CASE: $x, y \in \alpha$
- PROOF: The result follows because α is well-ordered by \in .
- $\langle 3 \rangle 2$. CASE: $x \in \alpha, y = \alpha$
- PROOF: Then $x \in y$.
- $\langle 3 \rangle 3$. CASE: $x = \alpha, y \in \alpha$
- PROOF: Then $y \in x$.
- $\langle 3 \rangle 4$. CASE: $x = \alpha, y = \alpha$
- PROOF: Then $x = y$.
- $\langle 2 \rangle 3$. Every nonempty subset of α^+ has an \in -least element.
- $\langle 3 \rangle 1$. LET: $S \subseteq \alpha^+$ be nonempty
- $\langle 3 \rangle 2$. CASE: $S = \{\alpha\}$
- PROOF: α is least in S .
- $\langle 3 \rangle 3$. CASE: $S \neq \{\alpha\}$
- $\langle 4 \rangle 1$. $S - \{\alpha\}$ is a nonempty subset of α
- $\langle 4 \rangle 2$. LET: β be least in $S - \{\alpha\}$
- $\langle 4 \rangle 3$. β is least in S .

□

Proposition 6.2.7. For ordinals α and β , if $\alpha^+ = \beta^+$ then $\alpha = \beta$.

PROOF: If $\alpha^+ = \beta^+$ then

$$\begin{aligned}\alpha &= \bigcup (\alpha^+) && \text{(Proposition 6.2.2)} \\ &= \bigcup (\beta^+) \\ &= \beta && \text{(Proposition 6.2.2)}\end{aligned}$$

Proposition 6.2.8. *For ordinals α and β , we have $\alpha < \beta$ if and only if $\alpha^+ < \beta^+$.*

PROOF:

$$\begin{aligned}\alpha < \beta &\Leftrightarrow \alpha^+ \leq \beta \\ &\Leftrightarrow \alpha^+ < \beta^+ && \square\end{aligned}$$

Definition 6.2.9 (Successor Ordinal). An ordinal α is a *successor ordinal* iff $\alpha = \beta^+$ for some β .

Definition 6.2.10 (Limit Ordinal). A *limit ordinal* is an ordinal that is neither 0 nor a successor ordinal.

Proposition 6.2.11. *If λ is a limit ordinal and $\beta < \lambda$ then $\beta^+ < \lambda$.*

PROOF: Since $\beta^+ \leq \lambda$ and $\beta^+ \neq \lambda$. \square

6.3 The Well-Ordering Theorem and Zorn's Lemma

Theorem 6.3.1 (Hartogs). *For any set A , there exists an ordinal not dominated by A .*

PROOF:

$\langle 1 \rangle 1$. LET: α be the class of all ordinals β such that $\beta \prec A$

PROVE: α is a set.

$\langle 1 \rangle 2$. LET: $W = \{(B, R) \mid B \subseteq A, R \text{ is a well ordering on } B\}$

$\langle 1 \rangle 3$. α is the class of the ordinals of the elements of W .

$\langle 2 \rangle 1$. For all $(B, R) \in W$, the ordinal of (B, R) is in α .

$\langle 3 \rangle 1$. LET: $(B, R) \in W$

$\langle 3 \rangle 2$. LET: β be the ordinal of (B, R)

$\langle 3 \rangle 3$. LET: $E : B \cong \beta$ be the canonical isomorphism.

$\langle 3 \rangle 4$. LET: $i : B \hookrightarrow A$ be the inclusion

$\langle 3 \rangle 5$. $i \circ E^{-1}$ is an injection $\beta \rightarrow A$

$\langle 3 \rangle 6$. $\beta \in \alpha$

$\langle 2 \rangle 2$. For all $\beta \in \alpha$, there exists $(B, R) \in W$ such that β is the ordinal number of (B, R) .

$\langle 3 \rangle 1$. LET: $\beta \in \alpha$

$\langle 3 \rangle 2$. PICK an injection $f : \beta \rightarrow A$

$\langle 3 \rangle 3$. Define \leq on $\text{ran } f$ by $f(x) \leq f(y)$ iff $x \leq y$

$\langle 3 \rangle 4$. $(\text{ran } f, \leq) \in W$

$\langle 3 \rangle 5$. β is the ordinal number of $(\text{ran } f, \leq)$

⟨1⟩4. α is a set.

PROOF: By an Axiom of Replacement.

⟨1⟩5. α is an ordinal.

PROOF: It is a transitive set of ordinals.

⟨1⟩6. $\alpha \not\subseteq A$

PROOF: Since $\alpha \notin \alpha$.

□

Theorem 6.3.2 (Numeration Theorem). *Every set is equinumerous with some ordinal.*

PROOF:

⟨1⟩1. LET: A be any set.

⟨1⟩2. PICK an ordinal α not dominated by A .

⟨1⟩3. PICK a choice function G for A .

⟨1⟩4. PICK $e \notin A$

⟨1⟩5. LET: $F : \alpha \rightarrow A \cup \{e\}$ by transfinite recursion:

$$F(\gamma) = \begin{cases} G(A - F(\{\delta \mid \delta < \gamma\})) & \text{if } A - F(\{\delta \mid \delta < \gamma\}) \neq \emptyset \\ e & \text{if } A - F(\{\delta \mid \delta < \gamma\}) = \emptyset \end{cases}$$

⟨1⟩6. $e \in \text{ran } F$

⟨2⟩1. ASSUME: for a contradiction $e \notin \text{ran } F$

⟨2⟩2. F is an injection $\alpha \rightarrow A$.

⟨3⟩1. LET: $\beta, \gamma \in \alpha$ with $\beta \neq \gamma$

PROVE: $F(\beta) \neq F(\gamma)$

⟨3⟩2. ASSUME: w.l.o.g. $\beta < \gamma$

⟨3⟩3. $F(\gamma) \in A - F(\{\delta \mid \delta < \gamma\})$

⟨3⟩4. $F(\gamma) \notin F(\{\delta \mid \delta < \gamma\})$

⟨3⟩5. $F(\gamma) \neq F(\beta)$

⟨2⟩3. Q.E.D.

PROOF: This contradicts ⟨1⟩2.

⟨1⟩7. LET: δ be least such that $F(\delta) = e$

⟨1⟩8. $F \upharpoonright \delta : \delta \approx A$

Theorem 6.3.3 (Well-Ordering Theorem). *Any set can be well ordered.*

PROOF:

⟨1⟩1. PICK an ordinal δ and a bijection $F : A \approx \delta$

⟨1⟩2. Define \leq on A by $F(x) \leq F(y)$ iff $x \leq y$ for $x, y \in \delta$

⟨1⟩3. \leq is a well ordering on A .

□

Theorem 6.3.4 (Zorn's Lemma). *Let \mathcal{A} be a set such that, for every chain $\mathcal{B} \subseteq \mathcal{A}$, we have $\bigcup \mathcal{B} \in \mathcal{A}$. Then \mathcal{A} has a maximal element.*

PROOF:

⟨1⟩1. PICK a well ordering $<$ on \mathcal{A} .

- ⟨1⟩2. LET: $F : \mathcal{A} \rightarrow 2$ be the function defined by transfinite recursion by:
- $$F(A) = \begin{cases} 1 & \text{if } A \text{ includes every set } B < A \text{ for which } F(B) = 1 \\ 0 & \text{otherwise} \end{cases}$$
- ⟨1⟩3. LET: $\mathcal{C} = \{A \in \mathcal{A} \mid F(A) = 1\}$
 PROVE: $\bigcup \mathcal{C}$ is a maximal element of \mathcal{A}
- ⟨1⟩4. For all $A \in \mathcal{A}$, we have $A \in \mathcal{C}$ iff $\forall B < A. B \in \mathcal{C} \Rightarrow B \subseteq A$
- ⟨1⟩5. \mathcal{C} is a chain.
- ⟨2⟩1. LET: $A, A' \in \mathcal{C}$
- ⟨2⟩2. ASSUME: w.l.o.g. $A \leq A'$
- ⟨2⟩3. $A \subseteq A'$
- PROOF: By ⟨1⟩4
- ⟨1⟩6. $\bigcup \mathcal{C} \in \mathcal{A}$
- ⟨1⟩7. $\bigcup \mathcal{C}$ is maximal in \mathcal{A} .
- ⟨2⟩1. LET: $A \in \mathcal{A}$ and $\bigcup \mathcal{C} \subseteq A$
- ⟨2⟩2. $A \in \mathcal{C}$
- PROOF: By ⟨1⟩4 since $\forall B \in \mathcal{C}. B \subseteq A$.
- ⟨2⟩3. $A \subseteq \bigcup \mathcal{C}$
- ⟨2⟩4. $A = \bigcup \mathcal{C}$
-

Proposition 6.3.5 (Teichmüller-Tukey Lemma). *Let \mathcal{A} be a nonempty set such that, for every B , we have $B \in \mathcal{A}$ if and only if every finite subset of B is a member of \mathcal{A} . Then \mathcal{A} has a maximal element.*

PROOF:

- ⟨1⟩1. For every chain $\mathcal{B} \subseteq \mathcal{A}$, we have $\bigcup \mathcal{B} \in \mathcal{A}$
- ⟨2⟩1. LET: $\mathcal{B} \subseteq \mathcal{A}$ be a chain.
- ⟨2⟩2. Every finite subset of $\bigcup \mathcal{B}$ is a member of \mathcal{A} .
- ⟨3⟩1. LET: C be a finite subset of $\bigcup \mathcal{B}$.
- ⟨3⟩2. PICK $B \in \mathcal{B}$ such that $C \subseteq B$.
- ⟨3⟩3. $B \in \mathcal{A}$
- ⟨3⟩4. Every finite subset of B is in \mathcal{A} .
- ⟨3⟩5. $C \in \mathcal{A}$
- ⟨2⟩3. $\bigcup \mathcal{B} \in \mathcal{A}$.
- ⟨1⟩2. Q.E.D.
- PROOF: Zorn's lemma.
-

Theorem Schema 6.3.6. *For any class \mathbf{A} , there exists a class \mathbf{F} such that the following is a theorem:*

If \mathbf{A} is a proper class of ordinals, then $\mathbf{F} : \mathbf{On} \rightarrow \mathbf{A}$ is an order isomorphism.

PROOF:

- ⟨1⟩1. Define $\mathbf{F} : \mathbf{On} \rightarrow \mathbf{A}$ by transfinite recursion as follows: $\mathbf{F}(\alpha)$ is the least element of \mathbf{A} that is different from $\mathbf{F}(\beta)$ for all $\beta < \alpha$.
- ⟨1⟩2. For all $\alpha, \beta \in \mathbf{On}$, if $\alpha < \beta$ then $\mathbf{F}(\alpha) < \mathbf{F}(\beta)$

PROOF: We have $\mathbf{F}(\alpha) \neq \mathbf{F}(\beta)$ by the definition of $\mathbf{F}(\beta)$, and $\mathbf{F}(\beta) \not\prec \mathbf{F}(\alpha)$ by the leastness of $\mathbf{F}(\alpha)$.

$\langle 1 \rangle 3$. \mathbf{F} is surjective.

$\langle 2 \rangle 1$. LET: $\alpha \in \mathbf{A}$

$\langle 2 \rangle 2$. ASSUME: as transfinite induction hypothesis $\forall \beta \in \mathbf{A}$, if $\beta < \alpha$ then there exists γ such that $\beta = \mathbf{F}(\gamma)$.

$\langle 2 \rangle 3$. LET: $\gamma = \{\delta \in \mathbf{On} \mid \mathbf{F}(\delta) < \alpha\}$

$\langle 2 \rangle 4$. γ is a set.

PROOF: Axiom of Replacement applied to α .

$\langle 2 \rangle 5$. γ is a transitive set.

PROOF: If $\mathbf{F}(\delta) < \alpha$ and $\epsilon < \delta$ then $\mathbf{F}(\epsilon) < \alpha$ by $\langle 1 \rangle 2$.

$\langle 2 \rangle 6$. γ is an ordinal.

PROOF: Proposition 6.1.8.

$\langle 2 \rangle 7$. $\mathbf{F}(\gamma) = \alpha$

$\langle 3 \rangle 1$. $\mathbf{F}(\gamma)$ is the least element of \mathbf{A} different from $\mathbf{F}(\delta)$ for all $\delta < \gamma$

$\langle 3 \rangle 2$. $\mathbf{F}(\gamma)$ is the least element of \mathbf{A} different from x for all $x \in \mathbf{A}$ with $x < \alpha$

$\langle 3 \rangle 3$. $\mathbf{F}(\gamma) = \alpha$

□

6.4 Ordinal Operations

Definition 6.4.1 (Ordinal Operation). An *ordinal operation* is a function $\mathbf{On} \rightarrow \mathbf{On}$.

Definition 6.4.2 (Continuous). An ordinal operation $\mathbf{T} : \mathbf{On} \rightarrow \mathbf{On}$ is *continuous* iff, for every limit ordinal λ , we have $\mathbf{T}(\lambda) = \bigcup_{\alpha < \lambda} \mathbf{T}(\alpha)$.

Definition 6.4.3 (Normal). An ordinal operation is *normal* iff it is continuous and strictly monotone.

Proposition Schema 6.4.4. For any class \mathbf{T} , the following is a theorem.

If \mathbf{T} is a continuous ordinal operation and $\forall \gamma. \mathbf{T}(\gamma) < \mathbf{T}(\gamma^+)$, then \mathbf{T} is normal.

PROOF:

$\langle 1 \rangle 1$. LET: $P[\beta]$ be the property $\forall \gamma < \beta. \mathbf{T}(\gamma) < \mathbf{T}(\beta)$

$\langle 1 \rangle 2$. $P[0]$

PROOF: Vacuous.

$\langle 1 \rangle 3$. For any ordinal γ , if $P[\gamma]$ then $P[\gamma^+]$

$\langle 2 \rangle 1$. ASSUME: $P[\gamma]$

$\langle 2 \rangle 2$. LET: $\delta < \gamma^+$

$\langle 2 \rangle 3$. CASE: $\delta < \gamma$

PROOF: Then $\mathbf{T}(\delta) < \mathbf{T}(\gamma) < \mathbf{T}(\gamma^+)$.

$\langle 2 \rangle 4$. CASE: $\delta = \gamma$

PROOF: Then $\mathbf{T}(\delta) = \mathbf{T}(\gamma) < \mathbf{T}(\gamma^+)$.

$\langle 1 \rangle 4$. For any limit ordinal λ , if $\forall \gamma < \lambda. P[\gamma]$ then $P[\lambda]$.

$\langle 2 \rangle 1$. ASSUME: $\forall \gamma < \lambda. P[\gamma]$

$\langle 2 \rangle 2$. LET: $\delta < \lambda$

$\langle 2 \rangle 3$. $\mathbf{T}(\delta) < \mathbf{T}(\lambda)$

PROOF:

$$\begin{aligned} \mathbf{T}(\delta) &< \mathbf{T}(\delta^+) \\ &\leq \bigcup_{\epsilon < \lambda} \mathbf{T}(\epsilon) \\ &= \mathbf{T}(\lambda) \end{aligned}$$

□

Proposition Schema 6.4.5. *For any class \mathbf{T} , the following is a theorem:*

Assume \mathbf{T} is a normal ordinal operation. For every ordinal α , we have $\alpha \leq \mathbf{T}(\alpha)$.

PROOF:

$\langle 1 \rangle 1$. LET: γ be an ordinal.

$\langle 1 \rangle 2$. ASSUME: as induction hypothesis $\forall \delta < \gamma. \mathbf{T}(\delta) \geq \delta$

$\langle 1 \rangle 3$. For all $\delta < \gamma$ we have $\delta < \mathbf{T}(\gamma)$

PROOF: \mathbf{T} is strictly monotone.

$\langle 1 \rangle 4$. $\gamma \leq \mathbf{T}(\gamma)$

□

Proposition Schema 6.4.6. *For any class \mathbf{T} , the following is a theorem:*

Assume \mathbf{T} is a normal ordinal operation. For any ordinal $\beta \geq \mathbf{T}(0)$, there exists a greatest ordinal γ such that $\mathbf{T}(\gamma) \leq \beta$.

PROOF:

$\langle 1 \rangle 1$. There exists γ such that $\mathbf{T}(\gamma) > \beta$

$\langle 2 \rangle 1$. For all γ we have $\mathbf{T}(\gamma) \geq \gamma$

PROOF: Proposition 6.4.5.

$\langle 2 \rangle 2$. $\mathbf{T}(\beta^+) > \beta$

$\langle 1 \rangle 2$. LET: δ be least such that $\mathbf{T}(\delta) > \beta$

$\langle 1 \rangle 3$. δ is a successor ordinal.

$\langle 2 \rangle 1$. $\delta \neq 0$

PROOF: Since $\mathbf{T}(0) \leq \beta$.

$\langle 2 \rangle 2$. δ is not a limit ordinal.

$\langle 3 \rangle 1$. ASSUME: for a contradiction δ is a limit ordinal.

$\langle 3 \rangle 2$. $\beta < \bigcup_{\epsilon < \delta} \mathbf{T}(\epsilon)$

PROOF: \mathbf{T} is continuous.

$\langle 3 \rangle 3$. There exists $\epsilon < \delta$ such that $\beta < \mathbf{T}(\epsilon)$

$\langle 3 \rangle 4$. Q.E.D.

PROOF: This contradicts the minimality of δ .

$\langle 1 \rangle 4$. LET: $\delta = \gamma^+$

$\langle 1 \rangle 5$. γ is greatest such that $\mathbf{T}(\gamma) \leq \beta$

□

Theorem Schema 6.4.7. *For any class \mathbf{T} , the following is a theorem:*

Assume that \mathbf{T} is a normal ordinal operation. For any nonempty set of ordinals S , we have

$$\mathbf{T}(\sup S) = \sup_{\alpha \in S} \mathbf{T}(\alpha) .$$

PROOF:

$\langle 1 \rangle 1.$ $\forall \alpha \in S. \mathbf{T}(\alpha) \leq \mathbf{T}(\sup S)$

PROOF: Since \mathbf{T} is monotone.

$\langle 1 \rangle 2.$ For any ordinal β , if $\forall \alpha \in S. \mathbf{T}(\alpha) \leq \beta$, then $\mathbf{T}(\sup S) \leq \beta$

$\langle 2 \rangle 1.$ LET: β be an ordinal.

$\langle 2 \rangle 2.$ LET: $\gamma = \sup S$

$\langle 2 \rangle 3.$ ASSUME: $\forall \alpha \in S. \mathbf{T}(\alpha) \leq \beta$

$\langle 2 \rangle 4.$ CASE: γ is 0 or a successor ordinal

PROOF: Then we must have $\gamma \in S$ so $\mathbf{T}(\gamma) \leq \beta$ from $\langle 2 \rangle 3.$

$\langle 2 \rangle 5.$ CASE: γ is a limit ordinal

$\langle 3 \rangle 1.$ $\mathbf{T}(\gamma) = \sup_{\alpha < \gamma} \mathbf{T}(\alpha)$

PROOF: \mathbf{T} is continuous.

$\langle 3 \rangle 2.$ ASSUME: for a contradiction $\beta < \mathbf{T}(\gamma)$

$\langle 3 \rangle 3.$ PICK $\alpha < \gamma$ such that $\beta < \mathbf{T}(\alpha)$

PROOF: $\langle 3 \rangle 1, \langle 3 \rangle 2$

$\langle 3 \rangle 4.$ PICK $\alpha' \in S$ such that $\alpha < \alpha'$

PROOF: $\langle 2 \rangle 2, \langle 3 \rangle 3$

$\langle 3 \rangle 5.$ $\beta < \mathbf{T}(\alpha') \leq \beta$

PROOF: \mathbf{T} is strictly monotone, $\langle 3 \rangle 3, \langle 3 \rangle 4, \langle 2 \rangle 3.$

$\langle 3 \rangle 6.$ Q.E.D.

PROOF: This is a contradiction.

□

Proposition 6.4.8. For any classes \mathbf{A} and \mathbf{T} , the following is a theorem:

Assume \mathbf{A} is a proper class of ordinals such that, for every set $S \subseteq \mathbf{A}$, we have $\bigcup S \in \mathbf{A}$. Assume \mathbf{T} is the unique order isomorphism $\mathbf{On} \cong \mathbf{A}$. Then \mathbf{T} is normal.

PROOF:

$\langle 1 \rangle 1.$ \mathbf{T} is strictly monotone.

PROOF: Since it is an order isomorphism.

$\langle 1 \rangle 2.$ \mathbf{T} is continuous.

$\langle 2 \rangle 1.$ LET: λ be a limit ordinal.

$\langle 2 \rangle 2.$ $\mathbf{T}'(\lambda)$ is the least member of \mathbf{A} that is greater than $\mathbf{T}'(\alpha)$ for all $\alpha < \lambda$

$\langle 2 \rangle 3.$ $\mathbf{T}'(\lambda) = \sup_{\alpha < \lambda} \mathbf{T}'(\alpha)$

□

Proposition Schema 6.4.9. For any class \mathbf{T} , the following is a theorem:

If \mathbf{T} is a normal ordinal operation, then for any limit ordinal λ , we have $\mathbf{T}(\lambda)$ is a limit ordinal.

PROOF:

$\langle 1 \rangle 1.$ $\mathbf{T}(\lambda) \neq 0$

PROOF: Since $0 \leq \mathbf{T}(0) < \mathbf{T}(\lambda)$.

$\langle 1 \rangle 2$. $\mathbf{T}(\lambda)$ is not a successor ordinal.

$\langle 2 \rangle 1$. ASSUME: for a contradiction $\mathbf{T}(\lambda) = \alpha^+$

$\langle 2 \rangle 2$. $\alpha < \mathbf{T}(\lambda) = \sup_{\beta < \lambda} \mathbf{T}(\beta)$

$\langle 2 \rangle 3$. PICK $\beta < \lambda$ such that $\alpha < \mathbf{T}(\beta)$

$\langle 2 \rangle 4$. $\alpha^+ \leq \mathbf{T}(\beta) < \mathbf{T}(\lambda)$

$\langle 2 \rangle 5$. Q.E.D.

PROOF: This is a contradiction.

□

6.5 Ordinal Arithmetic

6.5.1 Addition

Definition 6.5.1. Let A and B be disjoint well-ordered sets. The *concatenation* of A and B is the set $A \cup B$ under the relation:

- if $a, a' \in A$ then $a \leq a'$ iff $a \leq a'$ in A
- if $b, b' \in B$ then $b \leq b'$ iff $b \leq b'$ in B
- if $a \in A$ and $b \in B$ then $a \leq b$ and $b \not\leq a$.

Proposition 6.5.2. If A and B are disjoint well-ordered sets, then their concatenation is well-ordered.

PROOF:

$\langle 1 \rangle 1$. \leq is reflexive.

PROOF: For all $a \in A$ we have $a \leq a$, and for all $b \in B$ we have $b \leq b$.

$\langle 1 \rangle 2$. \leq is antisymmetric.

$\langle 2 \rangle 1$. ASSUME: $x \leq y \leq x$

$\langle 2 \rangle 2$. CASE: $x, y \in A$

PROOF: Then $x = y$ since the order on A is antisymmetric.

$\langle 2 \rangle 3$. CASE: $x \in A$ and $y \in B$

PROOF: This is impossible as it would imply $y \leq x$.

$\langle 2 \rangle 4$. CASE: $x \in B$ and $y \in A$

PROOF: This is impossible as it would imply $x \leq y$.

$\langle 2 \rangle 5$. CASE: $x, y \in B$

PROOF: Then $x = y$ since the order on B is antisymmetric.

$\langle 1 \rangle 3$. \leq is transitive.

$\langle 2 \rangle 1$. ASSUME: $x \leq y \leq z$

$\langle 2 \rangle 2$. CASE: $x, z \in A$

PROOF: In this case $y \in A$ since $y \leq z$, and so $x \leq z$ since the order on A is transitive.

$\langle 2 \rangle 3$. CASE: $x \in A$ and $z \in B$

PROOF: Then $x \leq z$ immediately.

$\langle 2 \rangle 4$. CASE: $x \in B$ and $z \in A$

PROOF: This is impossible because we have $y \notin A$ since $x \leq y$ and $y \notin B$ since $y \leq z$.

$\langle 2 \rangle 5$. CASE: $x, z \in B$

PROOF: In this case $y \in B$ since $x \leq y$, and so $x \leq z$ since the order on B is transitive.

$\langle 1 \rangle 4$. \leq is total.

$\langle 2 \rangle 1$. LET: $x, y \in A \cup B$

$\langle 2 \rangle 2$. CASE: $x, y \in A$

PROOF: Then $x \leq y$ or $y \leq x$ because the order on A is total.

$\langle 2 \rangle 3$. CASE: $x \in A$ and $y \in B$

PROOF: Then $x \leq y$.

$\langle 2 \rangle 4$. CASE: $x \in B$ and $y \in A$

PROOF: Then $y \leq x$.

$\langle 2 \rangle 5$. CASE: $x, y \in B$

PROOF: Then $x \leq y$ or $y \leq x$ because the order on B is total.

$\langle 1 \rangle 5$. Every nonempty subset of $A \cup B$ has a least element.

$\langle 2 \rangle 1$. LET: S be a nonempty subset of $A \cup B$

$\langle 2 \rangle 2$. CASE: $S \cap A = \emptyset$

PROOF: Then $S \subseteq B$ and so S has a least element.

$\langle 2 \rangle 3$. CASE: $S \cap A \neq \emptyset$

PROOF: The least element of $S \cap A$ is the least element of S .

□

Definition 6.5.3 (Ordinal Addition). Let α and β be ordinal numbers. Then $\alpha + \beta$ is the ordinal number of the concatenation of A and B , where A is any well ordered set with ordinal α and B is any well ordered set with ordinal β .

Theorem 6.5.4 (Associative Law for Addition). For any ordinals ρ , σ and τ , we have

$$\rho + (\sigma + \tau) = (\rho + \sigma) + \tau .$$

PROOF: Given disjoint well ordered sets A , B and C , the concatenation of A with (the concatenation of B and C) is the same as the concatenation of (the concatenation of A and B) and C . □

Theorem 6.5.5. For any ordinal ρ we have

$$\rho + 0 = 0 + \rho = \rho .$$

PROOF: For any well ordered set A , the concatenation of A with \emptyset is A , and the concatenation of \emptyset with A is A . □

Theorem 6.5.6. For any ordinal α we have $\alpha + 1 = \alpha^+$.

PROOF: Since α^+ is the concatenation of α and $\{\alpha\}$. □

Theorem 6.5.7. For any ordinal α , the operation that maps β to $\alpha + \beta$ is normal.

PROOF:

$\langle 1 \rangle 1$. For any limit ordinal λ , we have $\alpha + \lambda = \sup_{\beta < \lambda} (\alpha + \beta)$.

$\langle 2 \rangle 1$. LET: λ be a limit ordinal.

$\langle 2 \rangle 2$. $(\{0\} \times \alpha) \cup (\{1\} \times \lambda) = \bigcup_{\beta \in \lambda} ((\{0\} \times \alpha) \cup (\{1\} \times \beta))$, where the order on the right hand side is as in Lemma 6.1.15.

PROOF:

$$\begin{aligned} (\{0\} \times \alpha) \cup (\{1\} \times \lambda) &= (\{0\} \times \alpha) \cup (\{1\} \times \bigcup_{\beta < \lambda} \beta) \\ &= (\{0\} \times \alpha) \cup \bigcup_{\beta < \lambda} (\{1\} \times \beta) \\ &= \bigcup_{\beta < \lambda} ((\{0\} \times \alpha) \cup (\{1\} \times \beta)) \end{aligned}$$

$\langle 1 \rangle 2$. For any ordinal β we have $\alpha + \beta < \alpha + \beta^+$

PROOF: Since $\alpha + \beta^+ = \alpha + \beta + 1 = (\alpha + \beta)^+$.

□

Corollary 6.5.7.1. *For any ordinals α , β , γ , we have $\beta < \gamma$ if and only if $\alpha + \beta < \alpha + \gamma$.*

Corollary 6.5.7.2 (Left Cancellation for Addition). *For any ordinals α , β and γ , if $\alpha + \beta = \alpha + \gamma$ then $\beta = \gamma$.*

Theorem 6.5.8. *For any ordinals α , β , γ , if $\beta \leq \gamma$ then $\beta + \alpha \leq \gamma + \alpha$.*

PROOF: Transfinite induction on α . □

Theorem 6.5.9 (Subtraction Theorem). *Let α and β be ordinals with $\alpha \leq \beta$. Then there exists a unique ordinal δ such that $\alpha + \delta = \beta$.*

PROOF:

$\langle 1 \rangle 1$. For all ordinals α and β with $\alpha \leq \beta$, there exists δ such that $\alpha + \delta = \beta$

$\langle 2 \rangle 1$. LET: α and β be ordinals with $\alpha \leq \beta$

$\langle 2 \rangle 2$. LET: δ be the greatest ordinal such that $\alpha + \delta \leq \beta$

PROOF: Proposition 6.4.6.

$\langle 2 \rangle 3$. $\alpha + \delta = \beta$

PROOF: If $\alpha + \delta < \beta$ then $\alpha + \delta + 1 \leq \beta$ contradicting the greatestness of δ .

$\langle 1 \rangle 2$. Q.E.D.

PROOF: Uniqueness follows from the Left Cancellation Law.

□

6.5.2 Multiplication

Definition 6.5.10 (Ordinal Multiplication). Let α and β be ordinal numbers. Then $\alpha\beta$ is the ordinal number of $A \times B$ under the lexicographic order, where A is any well ordered set with ordinal α and B is any well ordered set with ordinal β .

This is well defined by Proposition 5.3.5.

Theorem 6.5.11 (Associative Law). *For any ordinals ρ , σ and τ , we have*

$$\rho(\sigma\tau) = (\rho\sigma)\tau .$$

PROOF: Let A , B and C be well ordered sets with ordinals ρ , σ and τ . Then both $\rho(\sigma\tau)$ and $(\rho\sigma)\tau$ are the ordinal of $A \times B \times C$ under $(a, b, c) \leq (a', b', c') \Leftrightarrow a \leq a' \vee (a = a' \wedge b \leq b') \vee (a = a' \wedge b = b' \wedge c \leq c')$. \square

Theorem 6.5.12 (Left Distributive Law). *For any ordinals ρ , σ and τ , we have*

$$\rho(\sigma + \tau) = \rho\sigma + \rho\tau$$

PROOF: Let A , B and C be well ordered sets with ordinals ρ , σ and τ and with $B \cap C = \emptyset$. Then both $\rho(\sigma + \tau)$ and $\rho\sigma + \rho\tau$ are the ordinal of $A \times (B \cup C)$ under the lexicographic ordering. \square

Theorem 6.5.13. *For any ordinal ρ we have $\rho 0 = 0\rho = 0$.*

PROOF: For any well ordered set A we have $A \times \emptyset = \emptyset \times A = \emptyset$. \square

Theorem 6.5.14. *For any ordinal ρ we have $\rho 1 = 1\rho = \rho$.*

PROOF: Easy. \square

Theorem 6.5.15. *For any ordinals ρ and σ , if $\rho\sigma = 0$ then $\rho = 0$ or $\sigma = 0$.*

PROOF: If $A \times B = \emptyset$ then $A = \emptyset$ or $B = \emptyset$. \square

Theorem 6.5.16. *For any non-zero ordinal α , the operation that maps β to $\alpha\beta$ is normal.*

PROOF:

$\langle 1 \rangle 1$. For any limit ordinal λ , we have $\alpha\lambda = \bigcup_{\beta < \lambda} \alpha\beta$

$\langle 2 \rangle 1$. LET: λ be a limit ordinal

$\langle 2 \rangle 2$. $\alpha \times \lambda = \bigcup_{\beta < \lambda} (\alpha \times \beta)$ as well-ordered sets

$\langle 1 \rangle 2$. For any ordinal β we have $\alpha\beta < \alpha\beta^+$

PROOF: $\alpha\beta^+ = \alpha\beta + \alpha > \alpha\beta$

\square

Corollary 6.5.16.1. *For any ordinals α , β , γ , if $\alpha \neq 0$ then $\beta < \gamma$ if and only if $\alpha\beta < \alpha\gamma$.*

Corollary 6.5.16.2 (Left Cancellation for Multiplication). *For any ordinals α , β , γ , if $\alpha \neq 0$ and $\alpha\beta = \alpha\gamma$ then $\beta = \gamma$.*

Theorem 6.5.17. *For any ordinals α , β and γ , if $\beta \leq \gamma$ then $\beta\alpha \leq \gamma\alpha$.*

PROOF: Transfinite induction on α . \square

Theorem 6.5.18 (Division Theorem). *Let α and δ be ordinal numbers with $\delta \neq 0$. Then there exist unique ordinals β and γ with $\gamma < \delta$ and*

$$\alpha = \delta\beta + \gamma .$$

PROOF:

(1)1. For any ordinal numbers α and δ with $\delta \neq 0$, there exist ordinals β and γ such that $\gamma < \delta$ and $\alpha = \delta\beta + \gamma$

(2)1. LET: α and δ be ordinals with $\delta \neq 0$

(2)2. LET: β be the greatest ordinal such that $\delta\beta \leq \alpha$

PROOF: Proposition 6.4.6.

(2)3. There exists an ordinal γ such that $\alpha = \delta\beta + \gamma$

PROOF: Subtraction Theorem

(1)2. For any ordinals $\delta, \beta, \beta', \gamma, \gamma'$, if $\delta\beta + \gamma = \delta\beta' + \gamma'$ and $\delta \neq 0$ and $\gamma, \gamma' < \delta$ then $\beta = \beta'$ and $\gamma = \gamma'$

(2)1. LET: $\delta, \beta, \beta', \gamma, \gamma'$ be ordinals.

(2)2. ASSUME: $\delta \neq 0$ and $\delta\beta + \gamma = \delta\beta' + \gamma'$

(2)3. $\beta = \beta'$

(3)1. $\beta \not\leq \beta'$

PROOF: If $\beta < \beta'$ then

$$\begin{aligned} \delta\beta' + \gamma' &\geq \delta\beta' \\ &\geq \delta(\beta + 1) \\ &= \delta\beta + \delta \\ &> \delta\beta + \gamma \end{aligned}$$

(3)2. $\beta' \not\leq \beta$

PROOF: Similar.

(2)4. $\gamma = \gamma'$

PROOF: By Cancellation.

□

6.5.3 Exponentiation

Definition 6.5.19. Given ordinals α and β , define the ordinal α^β as follows:

$$\begin{aligned} 0^\alpha &:= 0 & (\alpha > 0) \\ \alpha^0 &:= 1 \\ \alpha^{\beta^+} &:= \alpha^\beta \alpha & (\alpha > 0) \\ \alpha^\lambda &:= \sup_{\beta < \lambda} \alpha^\beta & (\alpha > 0, \lambda \text{ a limit ordinal}) \end{aligned}$$

Theorem 6.5.20. Let α be an ordinal ≥ 2 . The operation that maps β to α^β is normal.

PROOF:

(1)1. For λ a limit ordinal we have $\alpha^\lambda = \sup_{\beta < \lambda} \alpha^\beta$

PROOF: By definition.

(1)2. For any ordinal β we have $\alpha^\beta < \alpha^{\beta^+}$

PROOF: We have $\alpha^{\beta^+} = \alpha^\beta \alpha > \alpha^\beta$ by Theorem 6.5.16 since $\alpha > 1$ and $\alpha^\beta \neq 0$.

□

Corollary 6.5.20.1. *For any ordinals α, β, γ , if $\alpha \geq 2$ then $\beta < \gamma$ if and only if $\alpha^\beta < \alpha^\gamma$.*

Corollary 6.5.20.2 (Cancellation for Exponentiation). *For any ordinals α, β, γ , if $\alpha \geq 2$ and $\alpha^\beta = \alpha^\gamma$ then $\beta = \gamma$.*

Theorem 6.5.21. *For any ordinals α, β and γ , if $\beta \leq \gamma$ then $\beta^\alpha \leq \gamma^\alpha$.*

PROOF: Transfinite induction on α .

Theorem 6.5.22 (Logarithm Theorem). *Let α and β be ordinal numbers with $\alpha \neq 0$ and $\beta > 1$. Then there exist unique ordinals γ, δ and ρ such that*

$$\alpha = \beta^\gamma \delta + \rho, \quad 0 \neq \delta < \beta, \quad \rho < \beta^\gamma.$$

PROOF:

(1)1. For any ordinals α and β with $\alpha \neq 0$ and $\beta > 1$, there exist ordinals γ, δ, ρ such that

$$\alpha = \beta^\gamma \delta + \rho, \quad 0 \neq \delta < \beta, \quad \rho < \beta^\gamma.$$

(2)1. LET: α and β be ordinals with $\alpha \neq 0$ and $\beta > 1$.

(2)2. LET: γ be the greatest ordinal such that $\beta^\gamma \leq \alpha$.

PROOF: Proposition 6.4.6.

(2)3. LET: δ and ρ be the unique ordinals with $\rho < \beta^\gamma$ such that $\alpha = \beta^\gamma \delta + \rho$.

PROOF: By the Division Theorem.

(2)4. $\delta \neq 0$

PROOF: If $\delta = 0$ then $\alpha = \beta^\gamma 0 + \rho = \rho < \beta^\gamma \leq \alpha$ which is a contradiction.

(2)5. $\delta < \beta$

PROOF: If $\beta \leq \delta$ then $\alpha \geq \beta^\gamma \delta \geq \beta^\gamma \beta = \beta^{\gamma+1}$, contradicting the greatestness of γ .

(1)2. If $\beta^\gamma \delta + \rho = \beta^{\gamma'} \delta' + \rho'$ with $\beta > 1$, $0 \neq \delta < \beta$, $0 \neq \delta' < \beta$, $\rho < \beta^\gamma$ and $\rho' < \beta^{\gamma'}$, then $\gamma = \gamma'$, $\delta = \delta'$ and $\rho = \rho'$.

(2)1. LET: $\alpha = \beta^\gamma \delta + \rho = \beta^{\gamma'} \delta' + \rho'$

(2)2. $\beta^\gamma \leq \alpha < \beta^{\gamma+1}$

(2)3. $\beta^{\gamma'} \leq \alpha < \beta^{\gamma'+1}$

(2)4. $\beta^\gamma < \beta^{\gamma'+1}$ and $\beta^{\gamma'} < \beta^{\gamma+1}$

(2)5. $\gamma < \gamma' + 1$ and $\gamma' < \gamma + 1$

(2)6. $\gamma = \gamma'$

(2)7. $\delta = \delta'$ and $\rho = \rho'$

PROOF: By the Division Theorem.

□

Theorem 6.5.23. *For any ordinal numbers α, β, γ , we have*

$$\alpha^{\beta+\gamma} = \alpha^\beta \alpha^\gamma.$$

PROOF:

(1)1. LET: $P[\gamma]$ be the property: for any ordinals α and β we have $\alpha^{\beta+\gamma} = \alpha^\beta \alpha^\gamma$

(1)2. $P[0]$

PROOF:

$$\begin{aligned}\alpha^{\beta+0} &= \alpha^\beta \\ &= \alpha^\beta 1 \\ &= \alpha^\beta \alpha^0\end{aligned}$$

$\langle 1 \rangle 3$. For all γ , if $P[\gamma]$ then $P[\gamma + 1]$

PROOF:

$$\begin{aligned}\alpha^{\beta+\gamma+1} &= \alpha^{\beta+\gamma} \alpha \\ &= \alpha^\beta \alpha^\gamma \alpha && \text{(induction hypothesis)} \\ &= \alpha^\beta \alpha^{\gamma+1}\end{aligned}$$

$\langle 1 \rangle 4$. For any limit ordinal λ , if $\forall \gamma < \lambda. P[\gamma]$ then $P[\lambda]$.

$\langle 2 \rangle 1$. LET: λ be a limit ordinal

$\langle 2 \rangle 2$. ASSUME: $\forall \gamma < \lambda. P[\gamma]$

$\langle 2 \rangle 3$. LET: α and β be any ordinals.

$\langle 2 \rangle 4$. CASE: $\alpha = 0$

PROOF: We have $\alpha^{\beta+\lambda} = \alpha^\beta \alpha^\lambda = 0$.

$\langle 2 \rangle 5$. CASE: $\alpha = 1$

PROOF: We have $\alpha^{\beta+\lambda} = \alpha^\beta \alpha^\lambda = 1$.

$\langle 2 \rangle 6$. CASE: $\alpha > 1$

PROOF:

$$\begin{aligned}\alpha^{\beta+\lambda} &= \alpha^{\sup_{\gamma < \lambda} (\beta+\gamma)} \\ &= \sup_{\gamma < \lambda} \alpha^{\beta+\gamma} && \text{(Theorem 6.4.7)} \\ &= \sup_{\gamma < \lambda} \alpha^\beta \alpha^\gamma && (\langle 2 \rangle 2) \\ &= \alpha^\beta \sup_{\gamma < \lambda} \alpha^\gamma && \text{(Theorem 6.4.7)} \\ &= \alpha^\beta \alpha^\lambda\end{aligned}$$

□

Theorem 6.5.24. For any ordinal numbers α , β and γ , we have

$$(\alpha^\beta)^\gamma = \alpha^{\beta\gamma}.$$

PROOF:

$\langle 1 \rangle 1$. LET: $P[\gamma]$ be the property: For any ordinals α and β , we have $(\alpha^\beta)^\gamma = \alpha^{\beta\gamma}$

$\langle 1 \rangle 2$. $P[0]$

PROOF:

$$\begin{aligned}(\alpha^\beta)^0 &= 1 \\ &= \alpha^{\beta 0}\end{aligned}$$

$\langle 1 \rangle 3$. $\forall \gamma \in \mathbf{On}. P[\gamma] \Rightarrow P[\gamma + 1]$

PROOF:

$$\begin{aligned}
 (\alpha^\beta)^{\gamma+1} &= (\alpha^\beta)^\gamma \alpha^\beta \\
 &= \alpha^{\beta\gamma} \alpha^\beta \\
 &= \alpha^{\beta\gamma+\beta} \\
 &= \alpha^{\beta(\gamma+1)}
 \end{aligned}$$

$\langle 1 \rangle 4$. For any limit ordinal λ , if $\forall \gamma < \lambda. P[\gamma]$ then $P[\lambda]$.

$\langle 2 \rangle 1$. LET: λ be a limit ordinal.

$\langle 2 \rangle 2$. ASSUME: $\forall \gamma < \lambda. P[\gamma]$

$\langle 2 \rangle 3$. LET: α and β be any ordinals.

$\langle 2 \rangle 4$. CASE: $\alpha = 0$ and $\beta = 0$

PROOF:

$$\begin{aligned}
 (0^\beta)^\lambda &= 1^\lambda \\
 &= 1 \\
 &= 0^0 \\
 &= 0^{0\lambda}
 \end{aligned}$$

$\langle 2 \rangle 5$. CASE: $\alpha = 0$ and $\beta \neq 0$

PROOF: $(0^\beta)^\lambda = 0^{\beta\lambda} = 0$.

$\langle 2 \rangle 6$. CASE: $\alpha = 1$

PROOF: $(1^\beta)^\lambda = 1^{\beta\lambda} = 1$

$\langle 2 \rangle 7$. CASE: $\alpha > 1$

PROOF:

$$\begin{aligned}
 (\alpha^\beta)^\lambda &= \sup_{\gamma < \lambda} (\alpha^\beta)^\gamma \\
 &= \sup_{\gamma < \lambda} \alpha^{\beta\gamma} \\
 &= \alpha^{\sup_{\gamma < \lambda} \beta\gamma} \\
 &= \alpha^{\beta\lambda}
 \end{aligned}$$

□

6.6 Sequences

i

Definition 6.6.1 (Sequence). Given an ordinal α and class \mathbf{A} , an α -sequence in \mathbf{A} is a function $a : \alpha \rightarrow \mathbf{A}$. We write a_β for $a(\beta)$, and $(a_\beta)_{\beta < \alpha}$ for a .

Definition 6.6.2 (Strictly Increasing). A sequence (a_β) of ordinals is *strictly increasing* iff, whenever $\beta < \gamma$, then $a_\beta < a_\gamma$.

Definition 6.6.3 (Subsequence). Let $(a_\beta)_{\beta < \gamma}$ be a sequence in \mathbf{A} . A *subsequence* of (a_β) is a sequence of the form $(a_{\beta_\xi})_{\xi < \delta}$ where $(\beta_\xi)_{\xi < \delta}$ is a strictly increasing sequence in γ .

Definition 6.6.4 (Convergence). Let $(a_\beta)_{\beta < \gamma}$ be a sequence of ordinals and λ an ordinal. Then (a_β) *converges* to the *limit* λ iff $\lambda = \sup_{\beta < \gamma} a_\beta$.

Lemma 6.6.5. *Let $(a_\beta)_{\beta < \gamma}$ be a sequence of ordinals. Then there is a strictly increasing subsequence $(a_{\beta_\xi})_{\xi < \delta}$ such that $\sup_{\xi < \delta} a_{\beta_\xi} = \sup_{\beta < \gamma} a_\beta$.*

PROOF: Define β_ξ by transfinite recursion as follows. β_ξ is the least β such that $a_\beta > a_{\beta_\zeta}$ for all $\zeta < \xi$ if there is such an a_β ; if not, the sequence ends. \square

6.7 Strict Supremum

Definition 6.7.1 (Strict Supremum). For any set S of ordinals, define the *strict supremum* of S , $\text{ssup } S$, to be the least ordinal greater than every member of S .

Chapter 7

Cardinal Numbers

7.1 Cardinal Numbers

Definition 7.1.1 (Cardinality). For any set A , the *cardinality* or *cardinal number* $|A|$ of A is the least ordinal equinumerous with A .

Let **Card** be the class of all cardinal numbers.

Proposition 7.1.2. For any sets A and B , we have $A \approx B$ iff $|A| = |B|$.

PROOF: Easy. \square

Definition 7.1.3 (Addition). Given cardinal numbers κ and λ , we define $\kappa + \lambda$ to be $|A \cup B|$ where A and B are disjoint sets of cardinality κ and λ respectively.

We prove this is well-defined.

PROOF:

$\langle 1 \rangle 1$. ASSUME: $A \approx A'$, $B \approx B'$, and $A \cap B = A' \cap B' = \emptyset$

$\langle 1 \rangle 2$. PICK bijections $f : A \approx A'$ and $g : B \approx B'$

$\langle 1 \rangle 3$. The function $A \cup B \rightarrow A' \cup B'$ that maps $a \in A$ to $f(a)$ and $b \in B$ to $g(b)$ is a bijection.

\square

Proposition 7.1.4. For any cardinal number κ , we have $\kappa + 0 = \kappa$.

PROOF: Let A and B be disjoint sets of cardinality κ and 0 . Then $B = \emptyset$ so $A \cup B = A$ and so $|A \cup B| = \kappa$. \square

Theorem 7.1.5 (Associative Law for Addition). For any cardinal numbers κ , λ , μ we have $\kappa + (\lambda + \mu) = (\kappa + \lambda) + \mu$.

PROOF: Since $A \cup (B \cup C) = (A \cup B) \cup C$. \square

Proposition 7.1.6. For any cardinal numbers κ and λ we have $\kappa + \lambda = \lambda + \kappa$.

PROOF: Since $A \cup B = B \cup A$. \square

Definition 7.1.7 (Multiplication). For κ and λ cardinal numbers, we define $\kappa\lambda$ to be the cardinal number of $A \times B$, where $|A| = \kappa$ and $|B| = \lambda$.

We prove this is well-defined.

PROOF: If $f : A \approx A'$ and $g : B \approx B'$ then the function that maps (a, b) to $(f(a), g(b))$ is a bijection $A \times B \approx A' \times B'$. \square

Proposition 7.1.8. For any cardinal number κ we have $\kappa \cdot 0 = 0$.

PROOF: Since $A \times \emptyset = \emptyset$. \square

Proposition 7.1.9. For any cardinal number κ we have $\kappa \cdot 1 = \kappa$.

PROOF: The function that maps (a, e) to a is a bijection $A \times \{e\} \approx A$. \square

Theorem 7.1.10 (Distributive Law). For any cardinal numbers κ , λ and μ , we have $\kappa(\lambda + \mu) = \kappa\lambda + \kappa\mu$.

PROOF: Since $A \times (B \cup C) = (A \times B) \cup (A \times C)$. \square

Theorem 7.1.11 (Associative Law for Multiplication). For any cardinal numbers κ , λ and μ , we have $\kappa(\lambda\mu) = (\kappa\lambda)\mu$.

PROOF: Since $A \times (B \times C) \approx (A \times B) \times C$. \square

Theorem 7.1.12 (Commutative Law for Multiplication). For any cardinal numbers κ and λ , we have $\kappa\lambda = \lambda\kappa$.

PROOF: Since $A \times B \approx B \times A$. \square

Theorem 7.1.13. For any cardinal numbers κ and λ , if $\kappa\lambda = 0$ then $\kappa = 0$ or $\lambda = 0$.

PROOF: if $A \times B = \emptyset$ then $A = \emptyset$ or $B = \emptyset$. \square

Definition 7.1.14 (Exponentiation). Given cardinal numbers κ and λ , we define κ^λ to be $|A^B|$, where $|A| = \kappa$ and $|B| = \lambda$.

We prove this is well-defined.

PROOF: If $f : A \approx A'$ and $g : B \approx B'$, then the function that maps $h : B \rightarrow A$ to $f \circ h \circ g^{-1}$ is a bijection $A^B \approx A'^{B'}$. \square

Proposition 7.1.15. For any cardinal numbers κ , λ and μ ,

$$\kappa^{\lambda+\mu} = (\kappa^\lambda)^\mu$$

PROOF: The function that maps $f : A \times B \rightarrow C$ to $\lambda a \in A. \lambda b \in B. f(a, b)$ is a bijection $A^{B \times C} \approx (A^B)^C$. \square

Proposition 7.1.16. For any cardinal numbers κ , λ and μ ,

$$(\kappa\lambda)^\mu = \kappa^\mu \lambda^\mu .$$

PROOF: The function $f : A^C \times B^C \rightarrow (A \times B)^C$ with $f(g, h)(c) = (g(c), h(c))$ is a bijection. \square

Proposition 7.1.17. *For any cardinal numbers κ , λ and μ , we have*

$$\kappa^{\lambda+\mu} = \kappa^\lambda \kappa^\mu .$$

PROOF: If $B \cap C = \emptyset$, then $f : A^B \times A^C \rightarrow A^{B \cup C}$ given by $f(g, h)(b) = g(b)$ and $f(g, h)(c) = h(c)$ is a bijection. \square

Proposition 7.1.18. *For any cardinal number κ , we have $\kappa^0 = 1$.*

PROOF: For any set A , we have $A^\emptyset = \{\emptyset\}$. \square

Proposition 7.1.19. *For any cardinal number κ , we have $\kappa^1 = \kappa$.*

PROOF: For any sets A and B , if $B = \{b\}$ then the function $f : A \rightarrow A^B$ with $f(a)(b) = a$ is a bijection. \square

Proposition 7.1.20. *For any non-zero cardinal number κ we have $0^\kappa = 0$.*

PROOF: If A is nonempty then there is no function $A \rightarrow \emptyset$. \square

Proposition 7.1.21. *For any set A we have $|\mathcal{P}A| = 2^{|A|}$.*

PROOF: The function $f : \mathcal{P}A \rightarrow 2^A$ where $f(X)(a) = 0$ if $a \notin X$ and $f(X)(a) = 1$ if $a \in X$. \square

Theorem 7.1.22 (König). *Let I be a set. Let $\{A_i\}_{i \in I}$ and $\{B_i\}_{i \in I}$ be families of sets. Assume that $\forall i \in I. |A_i| < |B_i|$. Then $|\bigcup_{i \in I} A_i| < |\prod_{i \in I} B_i|$.*

PROOF:

$\langle 1 \rangle 1$. For all $i \in I$, choose an injection $f_i : A_i \rightarrow B_i$

$\langle 1 \rangle 2$. For all $i \in I$, choose $b_i \in B_i - f_i(A_i)$

$\langle 1 \rangle 3$. $|\bigcup_{i \in I} A_i| \leq |\prod_{i \in I} B_i|$

$\langle 2 \rangle 1$. Define $g : \bigcup_{i \in I} A_i \rightarrow \prod_{i \in I} B_i$ by

$$g(i, a)(j) = \begin{cases} f_i(a) & \text{if } i = j \\ b_j & \text{otherwise} \end{cases}$$

$\langle 2 \rangle 2$. g is injective.

$\langle 1 \rangle 4$. $|\bigcup_{i \in I} A_i| \neq |\prod_{i \in I} B_i|$

$\langle 2 \rangle 1$. LET: $h : \bigcup_{i \in I} A_i \rightarrow \prod_{i \in I} B_i$

PROVE: h is not surjective.

$\langle 2 \rangle 2$. For $i \in I$, PICK $c_i \in B_i - \{h(i, a)(i) \mid i \in I\}$

$\langle 2 \rangle 3$. $c \in \prod_{i \in I} B_i$

$\langle 2 \rangle 4$. $c \notin \text{ran } h$

\square

Corollary 7.1.22.1. *For any cardinal number κ we have $\kappa < 2^\kappa$.*

7.2 Ordering on Cardinal Numbers

Definition 7.2.1. Given cardinal numbers κ and λ , we have $\kappa \leq \lambda$ iff $A \preccurlyeq B$, where $|A| = \kappa$ and $|B| = \lambda$.

PROOF:

$\langle 1 \rangle 1$. LET: $|A| = \kappa$ and $|B| = \lambda$

$\langle 1 \rangle 2$. PICK bijections $f : A \approx \kappa$ and $g : B \approx \lambda$

$\langle 1 \rangle 3$. If $\kappa \leq \lambda$ then $A \preccurlyeq B$

PROOF: Let $i : \kappa \hookrightarrow \lambda$ be the inclusion. Then $g^{-1} \circ i \circ f$ is an injection $A \rightarrow B$.

$\langle 1 \rangle 4$. If $A \preccurlyeq B$ then $\kappa \leq \lambda$

$\langle 2 \rangle 1$. ASSUME: $A \preccurlyeq B$

$\langle 2 \rangle 2$. PICK an injection $h : A \hookrightarrow B$

$\langle 2 \rangle 3$. $g(h(A)) \subseteq B$ is well-ordered by \in

$\langle 2 \rangle 4$. LET: γ be the ordinal number of $(g(h(A)), \in)$

$\langle 2 \rangle 5$. $\gamma \leq \lambda$

PROOF: Proposition 6.1.12.

$\langle 2 \rangle 6$. $\kappa \leq \gamma$

PROOF: By the leastness of κ , since A is equinumerous with γ .

$\langle 2 \rangle 7$. $\kappa \leq \lambda$

□

Corollary 7.2.1.1. *There is no largest cardinal number.*

Proposition 7.2.2. *For any cardinal numbers κ , λ , μ , if $\kappa \leq \lambda$ then $\kappa + \mu \leq \lambda + \mu$.*

PROOF: If $f : A \rightarrow B$ is injective, and $A \cap C = B \cap C = \emptyset$, then the function $A \cup C \rightarrow B \cup C$ that maps a to $f(a)$ and maps c to c is an injection. □

Proposition 7.2.3. *For any cardinal numbers κ , λ , μ , if $\kappa \leq \lambda$ then $\kappa\mu \leq \lambda\mu$.*

PROOF: If $f : A \rightarrow B$ is injective, then the function $A \times C \rightarrow B \times C$ that maps (a, c) to $(f(a), c)$ is injective. □

Proposition 7.2.4. *For any cardinal numbers κ , λ , μ , if $\kappa \leq \lambda$ then $\kappa^\mu \leq \lambda^\mu$.*

PROOF: Given an injection $f : A \rightarrow B$, the function that maps $A^C \rightarrow B^C$ that maps g to $f \circ g$ is an injection. □

Proposition 7.2.5. *For any cardinal numbers κ , λ , μ , if $\kappa \leq \lambda$ and μ and κ are not both 0, then $\mu^\kappa \leq \mu^\lambda$.*

PROOF:

$\langle 1 \rangle 1$. LET: A , B and C be sets with A and C not both empty.

$\langle 1 \rangle 2$. LET: $f : A \rightarrow B$ be an injection.

PROVE: $C^A \preccurlyeq C^B$

$\langle 1 \rangle 3$. CASE: $C = \emptyset$

PROOF: Then $A \neq \emptyset$ so $C^A = \emptyset \preccurlyeq C^B$.

$\langle 1 \rangle 4$. CASE: $C \neq \emptyset$

- ⟨2⟩1. PICK $c \in C$
 ⟨2⟩2. LET: $g : C^A \rightarrow C^B$ be the function $g(h)(f(a)) = h(a)$, $g(h)(b) = c$ if $b \notin f(A)$
 ⟨2⟩3. g is an injection.

□

Proposition 7.2.6. *Let \mathcal{A} be a set such that $\forall X \in \mathcal{A}. |X| \leq \kappa$. Then*

$$\left| \bigcup \mathcal{A} \right| \leq |\mathcal{A}| \kappa .$$

PROOF:

- ⟨1⟩1. For $X \in \mathcal{A}$, choose a surjection $f_X : \kappa \rightarrow X$.
 ⟨1⟩2. Define $g : \mathcal{A} \times \kappa \rightarrow \bigcup \mathcal{A}$ by $g(X, \alpha) = f_X(\alpha)$
 ⟨1⟩3. g is surjective.

□

Lemma 7.2.7. *The union of a set of cardinal numbers is a cardinal number.*

PROOF:

- ⟨1⟩1. LET: A be a set of cardinal numbers.
 PROVE: $\bigcup A$ is the smallest ordinal equinumerous with $\bigcup A$
 ⟨1⟩2. LET: $\alpha < \bigcup A$
 PROVE: $\alpha \not\approx \bigcup A$
 ⟨1⟩3. PICK $\kappa \in A$ such that $\alpha < \kappa$
 ⟨1⟩4. $\alpha \prec \kappa$
 ⟨1⟩5. $\alpha \prec \bigcup A$

□

Chapter 8

Natural Numbers

8.1 Inductive Sets

Definition 8.1.1 (Inductive). A set I is *inductive* iff $0 \in I$ and $\forall x \in I. x^+ \in I$.

Definition 8.1.2 (Natural Number). A *natural number* is a set that belongs to every inductive set.

Theorem 8.1.3. *The class \mathbb{N} of natural numbers is a set.*

PROOF:

$\langle 1 \rangle 1$. PICK an inductive set I .

PROOF: Axiom of Infinity.

$\langle 1 \rangle 2$. $\mathbb{N} \subseteq I$

□

Theorem 8.1.4. *\mathbb{N} is inductive, and is a subset of every other inductive set.*

PROOF:

$\langle 1 \rangle 1$. \mathbb{N} is inductive.

$\langle 2 \rangle 1$. $0 \in \mathbb{N}$

PROOF: Since 0 is a member of every inductive set.

$\langle 2 \rangle 2$. $\forall n \in \mathbb{N}. n^+ \in \mathbb{N}$

$\langle 3 \rangle 1$. LET: $n \in \mathbb{N}$

$\langle 3 \rangle 2$. LET: I be any inductive set.

PROVE: $n^+ \in I$

$\langle 3 \rangle 3$. $n \in I$

PROOF: $\langle 3 \rangle 1$, $\langle 3 \rangle 2$

$\langle 3 \rangle 4$. $n^+ \in I$

PROOF: $\langle 3 \rangle 2$, $\langle 3 \rangle 3$

$\langle 1 \rangle 2$. \mathbb{N} is a subset of every inductive set.

PROOF: Immediate from definitions.

□

Corollary 8.1.4.1 (Induction Principle for \mathbb{N}). *Any inductive subset of \mathbb{N} coincides with \mathbb{N} .*

Theorem 8.1.5. *Every natural number except 0 is the successor of some natural number.*

PROOF: Trivially by induction. \square

Proposition 8.1.6. *Every natural number is an ordinal.*

PROOF: By induction. \square

Proposition 8.1.7. *\mathbb{N} is a transitive set.*

PROOF:

$\langle 1 \rangle 1. 0 \subseteq \mathbb{N}$

$\langle 1 \rangle 2. \forall n \in \mathbb{N}. n \subseteq \mathbb{N} \Rightarrow n^+ \subseteq \mathbb{N}$

$\langle 1 \rangle 3. \forall n \in \mathbb{N}. n \subseteq \mathbb{N}$

PROOF: From $\langle 1 \rangle 1$ and $\langle 1 \rangle 2$ by induction.

\square

Corollary 8.1.7.1. *\mathbb{N} is an ordinal.*

Definition 8.1.8. We define $\omega = \mathbb{N}$.

Proposition 8.1.9 (Dependent Choice). *Let A be a nonempty set and R a relation on A such that $\forall x \in A. \exists y \in A. (y, x) \in R$. Then there exists a function $f : \mathbb{N} \rightarrow A$ such that $\forall n \in \mathbb{N}. (f(n+1), f(n)) \in R$.*

PROOF:

$\langle 1 \rangle 1.$ PICK a choice function F for A .

$\langle 1 \rangle 2.$ PICK $a \in A$

$\langle 1 \rangle 3.$ Define $f : \mathbb{N} \rightarrow A$ by $f(0) = a$ and $f(n+1) = F(\{y \in A \mid (y, f(n)) \in R\})$.

\square

Theorem Schema 8.1.10. *For any classes \mathbf{A} and \mathbf{R} , the following is a theorem:*

Assume \mathbf{R} is a relation on \mathbf{A} and, for all $a \in \mathbf{A}$, the class $\{x \in \mathbf{A} \mid x\mathbf{R}a\}$ is a set. Then \mathbf{R} is well founded if and only if there does not exist a function $f : \mathbb{N} \rightarrow \mathbf{A}$ such that $\forall n \in \mathbb{N}. f(n+1)\mathbf{R}f(n)$.

PROOF:

$\langle 1 \rangle 1.$ If there exists a function $f : \mathbb{N} \rightarrow \mathbf{A}$ such that $\forall n \in \mathbb{N}. f(n+1)\mathbf{R}f(n)$ then \mathbf{R} is not well founded.

PROOF: $f(\mathbb{N})$ is a nonempty subset of \mathbf{A} with no \mathbf{R} -minimal element.

$\langle 1 \rangle 2.$ If \mathbf{R} is not well founded then there exists a function $f : \mathbb{N} \rightarrow \mathbf{A}$ such that $\forall n \in \mathbb{N}. f(n+1)\mathbf{R}f(n)$.

$\langle 2 \rangle 1.$ ASSUME: \mathbf{R} is not well founded.

$\langle 2 \rangle 2.$ PICK a nonempty subset $B \subseteq \mathbf{A}$ that has no \mathbf{R} -minimal element.

$\langle 2 \rangle 3.$ $\forall x \in B. \exists y \in B. y\mathbf{R}x$

- ⟨2⟩4. Choose a function $g : B \rightarrow B$ such that $\forall x \in B. g(x) \mathbf{R} x$
- ⟨2⟩5. PICK $b \in B$
- ⟨2⟩6. Define $f : \mathbb{N} \rightarrow \mathbf{A}$ recursively by $f(0) = b$ and $\forall n \in \mathbb{N}. f(n+1) = g(f(n))$
- ⟨2⟩7. $\forall n \in \mathbb{N}. f(n+1) \mathbf{R} f(n)$

□

8.2 Cardinality

Definition 8.2.1 (Finite). A set is *finite* iff it is equinumerous to some natural number; otherwise it is *infinite*.

Theorem 8.2.2 (Pigeonhole Principle). *No natural number is equinumerous to a proper subset of itself.*

PROOF:

⟨1⟩1. LET: $P(n)$ be the property: any one-to-one function $n \rightarrow n$ is surjective.

⟨1⟩2. $P(0)$

PROOF: The only function $0 \rightarrow 0$ is injective.

⟨1⟩3. For every natural number n , if $P(n)$ then $P(n+1)$.

⟨2⟩1. ASSUME: $P(n)$

⟨2⟩2. LET: f be a one-to-one function $n+1 \rightarrow n+1$

⟨2⟩3. $f \upharpoonright n$ is a one-to-one function $n \rightarrow n+1$

⟨2⟩4. CASE: $n \notin \text{ran } f$

⟨3⟩1. $f \upharpoonright n : n \rightarrow n$

⟨3⟩2. $\text{ran}(f \upharpoonright n) = n$

⟨3⟩3. $f(n) = n$

PROOF: ⟨2⟩1.

⟨3⟩4. $\text{ran } f = n+1$

⟨2⟩5. CASE: $n \in \text{ran } f$

⟨3⟩1. PICK $p \in n$ such that $f(p) = n$

⟨3⟩2. LET: $\hat{f} : n \rightarrow n$ be the function

$$\hat{f}(p) = f(n)$$

$$\hat{f}(x) = f(x) \quad (x \neq p)$$

⟨3⟩3. \hat{f} is one-to-one

⟨3⟩4. $\text{ran } \hat{f} = n$

PROOF: ⟨2⟩1

⟨3⟩5. $\text{ran } f = n+1$

⟨1⟩4. For every natural number n , $P(n)$.

□

Corollary 8.2.2.1. *No finite set is equinumerous to a proper subset of itself.*

Corollary 8.2.2.2. *Every natural number is a cardinal number.*

PROOF: For any natural number n , we have that n is the least ordinal such that $n \approx n$. □

Corollary 8.2.2.3. \mathbb{N} is a cardinal number.

Corollary 8.2.2.4. \mathbb{N} is infinite.

PROOF: The function that maps n to $n+1$ is a bijection between \mathbb{N} and $\mathbb{N}-\{0\}$.
 \square

Corollary 8.2.2.5. If C is a proper subset of a natural number n , then there exists $m < n$ such that $C \approx m$.

PROOF: By Proposition 6.1.12. \square

Corollary 8.2.2.6. Any subset of a finite set is finite.

Proposition 8.2.3. For any natural numbers m and n we have $m+n$ (cardinal addition) is a natural number.

PROOF: Induction on n . \square

Corollary 8.2.3.1. The union of two finite sets is finite.

Corollary 8.2.3.2. The union of a finite set of finite sets is finite.

PROOF: By induction on the number of elements. \square

Proposition 8.2.4. For natural numbers m and n , the cardinal sum $m+n$ is equal to the ordinal sum $m+n$.

PROOF: Induction on n . \square

Proposition 8.2.5. For any natural numbers m and n , we have mn (cardinal multiplication) is a natural number.

Corollary 8.2.5.1. If A and B are finite sets then $A \times B$ is finite.

Proposition 8.2.6. For natural numbers m and n , the cardinal product mn is equal to the ordinal product mn .

PROOF: Induction on n . \square

Proposition 8.2.7. For any natural numbers m and n we have m^n (cardinal exponentiation) is a natural number.

PROOF: Induction on n .

Corollary 8.2.7.1. If A and B are finite sets then A^B are finite.

Proposition 8.2.8. For natural numbers m and n , the cardinal exponentiation m^n and the ordinal exponentiation m^n agree.

PROOF: Induction on n . \square

Proposition 8.2.9. $\mathbb{N}^2 \approx \mathbb{N}$

PROOF: The function $J : \mathbb{N}^2 \rightarrow \mathbb{N}$ defined by $J(m, n) = ((m + n)^2 + 3m + n)/2$ is a bijection. \square

Proposition 8.2.10. *For any infinite cardinal κ we have $\aleph_0 \leq \kappa$.*

PROOF:

$\langle 1 \rangle 1$. LET: A be an infinite set.

PROVE: $\aleph_0 \leq A$

$\langle 1 \rangle 2$. PICK a choice function F for A .

$\langle 1 \rangle 3$. Define $h : \mathbb{N} \rightarrow \{X \in \mathcal{P}A \mid X \text{ is finite}\}$ by

$$h(0) = \emptyset$$

$$h(n+1) = h(n) \cup \{F(A - \{h(m) \mid m < n\})\}$$

$\langle 1 \rangle 4$. Define $g : \mathbb{N} \rightarrow A$ by $g(n) = F(A - \{h(m) \mid m < n\})$

$\langle 1 \rangle 5$. g is injective.

PROOF: If $m < n$ then $g(m) \neq g(n)$.

\square

Theorem Schema 8.2.11 (König's Lemma). *For any classes \mathbf{A} and \mathbf{R} , the following is a theorem:*

Assume \mathbf{R} is a well founded relation on \mathbf{A} such that, for all $y \in \mathbf{A}$, the class $\{x \in \mathbf{A} \mid x\mathbf{R}y\}$ is a finite set. Let \mathbf{R}^t be the transitive closure of \mathbf{R} . Then, for all $y \in \mathbf{A}$, the class $\{x \in \mathbf{A} \mid x\mathbf{R}^ty\}$ is a finite set.

PROOF:

$\langle 1 \rangle 1$. LET: $y \in \mathbf{A}$

$\langle 1 \rangle 2$. ASSUME: as transfinite induction hypothesis $\forall x\mathbf{R}y. \{z \in \mathbf{A} \mid z\mathbf{R}^tx\}$ is a finite set.

$\langle 1 \rangle 3$. $\{x \mid x\mathbf{R}^ty\} = \bigcup_{x\mathbf{R}y} (\{x\} \cup \{z \mid z\mathbf{R}^tx\})$

$\langle 1 \rangle 4$. $\{x \mid x\mathbf{R}^ty\}$ is finite.

PROOF: Corollary 8.2.3.2.

\square

8.3 Countable Sets

Definition 8.3.1 (Countable). A set A is *countable* iff $|A| \leq \aleph_0$.

Theorem 8.3.2. *The union of a countable set of countable sets is countable.*

PROOF: Proposition 7.2.6. \square

8.4 Arithmetic

Definition 8.4.1 (Even). A natural number n is *even* iff there exists $m \in \mathbb{N}$ such that $n = 2m$.

Definition 8.4.2 (Odd). A natural number n is *odd* iff there exists $p \in \mathbb{N}$ such that $n = 2p + 1$.

Proposition 8.4.3. *Every natural number is either even or odd.*

PROOF:

$\langle 1 \rangle 1$. 0 is even.

PROOF: $0 = 2 \times 0$.

$\langle 1 \rangle 2$. For any natural number n , if n is either even or odd then n^+ is either even or odd.

PROOF:

$\langle 2 \rangle 1$. LET: $n \in \mathbb{N}$

$\langle 2 \rangle 2$. If n is even then n^+ is odd.

PROOF: If $n = 2p$ then $n^+ = 2p + 1$.

$\langle 2 \rangle 3$. If n is odd then n^+ is even.

PROOF: If $n = 2p + 1$ then $n^+ = 2(p + 1)$.

□

Proposition 8.4.4. *No natural number is both even and odd.*

PROOF:

$\langle 1 \rangle 1$. 0 is not odd.

PROOF: For any p we have $2p + 1 = (2p)^+ \neq 0$.

$\langle 1 \rangle 2$. For any natural number n , if n is not both even and odd, then n^+ is not both even and odd.

$\langle 2 \rangle 1$. LET: n be a natural number.

$\langle 2 \rangle 2$. If n^+ is even then n is odd.

$\langle 3 \rangle 1$. ASSUME: n^+ is even.

$\langle 3 \rangle 2$. PICK p such that $n^+ = 2p$

$\langle 3 \rangle 3$. $p \neq 0$

PROOF: Since $n^+ \neq 0$.

$\langle 3 \rangle 4$. PICK q such that $p = q^+$

PROOF: Theorem 8.1.5.

$\langle 3 \rangle 5$. $n^+ = 2q + 2$

PROOF: $\langle 3 \rangle 2$, $\langle 3 \rangle 4$.

$\langle 3 \rangle 6$. $n = 2q + 1$

PROOF: Proposition 6.2.7, $\langle 3 \rangle 5$

$\langle 3 \rangle 7$. n is odd.

$\langle 2 \rangle 3$. If n^+ is odd then n is even.

$\langle 3 \rangle 1$. ASSUME: n^+ is odd.

$\langle 3 \rangle 2$. PICK p such that $n^+ = 2p + 1$

$\langle 3 \rangle 3$. $n = 2p$

PROOF: Proposition 6.2.7, $\langle 3 \rangle 2$

$\langle 3 \rangle 4$. n is even.

□

Proposition 8.4.5. *Let m, n, p, q be natural numbers. Assume $m + n = p + q$. Then $m < p$ if and only if $q < n$.*

PROOF:

$\langle 1 \rangle 1$. If $m < p$ then $q < n$.

PROOF: If $m < p$ and $n \leq q$ then $m + n < p + q$.

(1)2. If $q < n$ then $m < p$.

PROOF: Similar.

□

Proposition 8.4.6. *Let m, n, p and q be natural numbers. Assume $n < m$ and $q < p$. Then*

$$mq + np < mp + nq .$$

PROOF:

(1)1. PICK positive natural numbers a and b such that $m = n + a$ and $p = q + b$.

(1)2. $mp + nq > mq + np$

PROOF:

$$\begin{aligned} mp + nq &= (n + a)(q + b) + nq \\ &= 2nq + nb + aq + ab \\ mq + np &= (n + a)q + n(q + b) \\ &= 2nq + aq + nb \\ \therefore mp + nq &= mq + np + ab \\ &> mq + np \end{aligned}$$

□

8.5 Sequences

Definition 8.5.1 (Sequence). Let A be a set. A *finite sequence* in A is a function $a : n \rightarrow A$ for some natural number n ; we write it as $(a(0), a(1), \dots, a(n - 1))$. An (*infinite*) *sequence* in A is a function $\mathbb{N} \rightarrow A$.

We write A^* for the set of all finite sequences in A .

Proposition 8.5.2. *If A is countable then A^* is countable.*

PROOF: For any n , the set A^n is countable, and A^* is equinumerous with $\bigcup_n A^n$.

□

8.6 Transitive Closure of a Set

Proposition 8.6.1. *For any set A , there exists a unique transitive set C such that:*

- $A \subseteq C$
- For any transitive set X , if $A \subseteq X$ then $C \subseteq X$

PROOF:

(1)1. Define a function $F : \mathbb{N} \rightarrow \mathbf{V}$ by

$$F(0) = A$$

$$F(n + 1) = A \cup \bigcup (F(0) \cup \dots \cup F(n))$$

- $\langle 1 \rangle 2$. For all $n \in \mathbb{N}$ and $a \in F(n)$ we have $a \subseteq F(n+1)$
 PROOF: $a \in F(0) \cup \dots \cup F(n)$ so $a \subseteq \bigcup(F(0) \cup \dots \cup F(n)) \subseteq F(n+1)$.
 $\langle 1 \rangle 3$. LET: $C = \bigcup_{n \in \mathbb{N}} F(n)$
 $\langle 1 \rangle 4$. C is transitive.
 $\langle 2 \rangle 1$. LET: $x \in y \in C$
 $\langle 2 \rangle 2$. PICK $n \in \mathbb{N}$ such that $y \in F(n)$
 $\langle 2 \rangle 3$. $y \subseteq F(n+1)$
 PROOF: $\langle 1 \rangle 2$
 $\langle 2 \rangle 4$. $x \in F(n+1)$
 $\langle 2 \rangle 5$. $x \in C$
 $\langle 1 \rangle 5$. $A \subseteq C$
 PROOF: Since $F(0) = A$.
 $\langle 1 \rangle 6$. For any transitive set X , if $A \subseteq X$ then $C \subseteq X$
 $\langle 2 \rangle 1$. LET: X be a transitive set
 $\langle 2 \rangle 2$. ASSUME: $A \subseteq X$
 $\langle 2 \rangle 3$. For all $n \in \mathbb{N}$ we have $F(n) \subseteq X$.
 $\langle 3 \rangle 1$. $F(0) \subseteq X$
 PROOF: $\langle 2 \rangle 2$
 $\langle 3 \rangle 2$. For all $n \in \mathbb{N}$, if $F(n) \subseteq X$, then $F(n+1) \subseteq X$.
 $\langle 4 \rangle 1$. LET: $n \in \mathbb{N}$
 $\langle 4 \rangle 2$. ASSUME: $\forall m < n. F(m) \subseteq X$
 $\langle 4 \rangle 3$. $F(0) \cup \dots \cup F(n) \subseteq X$
 $\langle 4 \rangle 4$. $\bigcup(F(0) \cup \dots \cup F(n)) \subseteq X$
 PROOF: Since X is transitive.
 $\langle 4 \rangle 5$. $F(n+1) \subseteq X$
 $\langle 2 \rangle 4$. $C \subseteq X$
 $\langle 1 \rangle 7$. Let D be a transitive set such that $A \subseteq D$ and, for any transitive set X ,
 if $A \subseteq X$ then $D \subseteq X$. Then $D = C$.
 PROOF: We have $C \subseteq D$ and $D \subseteq C$.
 \square

8.7 The Veblen Fixed Point Theorem

Theorem Schema 8.7.1 (Veblen Fixed Point Theorem). *For any class \mathbf{T} , the following is a theorem:*

Assume \mathbf{T} is a normal ordinal operation. For every ordinal β , there exists $\gamma \geq \beta$ such that $\mathbf{T}(\gamma) = \gamma$.

PROOF:

- $\langle 1 \rangle 1$. LET: β be an ordinal.
 $\langle 1 \rangle 2$. ASSUME: w.l.o.g. $\beta < \mathbf{T}(\beta)$
 PROOF: We have $\beta \leq \mathbf{T}(\beta)$ by Proposition 6.4.5, and if $\beta = \mathbf{T}(\beta)$ we take $\gamma := \beta$.

$\langle 1 \rangle 3$. Define $f : \mathbb{N} \rightarrow \mathbf{On}$ by recursion thus:

$$\begin{aligned} f(0) &= \beta \\ f(n^+) &= \mathbf{T}(f(n)) \end{aligned}$$

$\langle 1 \rangle 4$. LET: $\gamma = \sup_{n \in \mathbb{N}} f(n)$

$\langle 1 \rangle 5$. $\beta \leq \gamma$

PROOF: Since $\beta = f(0)$.

$\langle 1 \rangle 6$. $\mathbf{T}(\gamma) = \gamma$

$\langle 2 \rangle 1$. $\mathbf{T}(\gamma) \leq \gamma$

PROOF:

$$\begin{aligned} \mathbf{T}(\gamma) &= \sup_{n \in \mathbb{N}} \mathbf{T}(f(n)) && (\text{Theorem 6.4.7}) \\ &= \sup_{n \in \mathbb{N}} f(n^+) && (\langle 1 \rangle 3) \\ &\leq \sup_{n \in \mathbb{N}} f(n) \\ &= \gamma \end{aligned}$$

$\langle 2 \rangle 2$. $\gamma \leq \mathbf{T}(\gamma)$

PROOF: Proposition 6.4.5.

□

Definition 8.7.2 (Derived Operation). Let \mathbf{T} be a normal ordinal operation. The *derived* operation $\mathbf{T}' : \mathbf{On} \rightarrow \mathbf{V}$ is the unique order isomorphism between \mathbf{On} and the fixed points of \mathbf{T} .

Proposition Schema 8.7.3. *For any class \mathbf{T} , the following is a theorem:*

If \mathbf{T} is a normal ordinal operation, then the derived operation is normal.

PROOF:

$\langle 1 \rangle 1$. For any set S of fixed points of \mathbf{T} , we have $\bigcup S$ is a fixed point of \mathbf{T}

$\langle 2 \rangle 1$. LET: S be a set of fixed points of \mathbf{T} .

$\langle 2 \rangle 2$. $\mathbf{T}(\sup S) = \sup S$

PROOF:

$$\begin{aligned} \mathbf{T}(\sup S) &= \sup_{\alpha \in S} \mathbf{T}(\alpha) && (\text{Theorem 6.4.7}) \\ &= \sup_{\alpha \in S} \alpha && (\langle 2 \rangle 1) \\ &= \sup S \end{aligned}$$

$\langle 1 \rangle 2$. Q.E.D.

PROOF: Proposition 6.4.8.

□

8.8 Cantor Normal Form

Theorem 8.8.1. *For any ordinal α , there exist a unique sequence of nonzero natural numbers (n_1, \dots, n_k) and sequence of ordinals $(\gamma_1, \dots, \gamma_k)$ such that*

$$\gamma_k < \gamma_{k-1} < \dots < \gamma_1$$

and

$$\alpha = \omega^{\gamma_1} n_1 + \omega^{\gamma_2} n_2 + \cdots + \omega^{\gamma_k} n_k .$$

PROOF:

$\langle 1 \rangle 1$. For any ordinal α , there exist a sequence of nonzero natural numbers (n_1, \dots, n_k) and sequence of ordinals $(\gamma_1, \dots, \gamma_k)$ such that

$$\gamma_k < \gamma_{k-1} < \cdots < \gamma_1$$

and

$$\alpha = \omega^{\gamma_1} n_1 + \omega^{\gamma_2} n_2 + \cdots + \omega^{\gamma_k} n_k .$$

$\langle 2 \rangle 1$. LET: α be an ordinal

$\langle 2 \rangle 2$. ASSUME: as an induction hypothesis that, for all $\beta < \alpha$, the theorem holds.

$\langle 2 \rangle 3$. ASSUME: w.l.o.g. $\alpha \neq 0$

$\langle 2 \rangle 4$. LET: γ_1, n_1, ρ_1 be the unique ordinals such that $0 \neq n_1 < \omega$, $\rho_1 < \omega^{\gamma_1}$, and $\alpha = \omega^{\gamma_1} n_1 + \rho_1$

$\langle 2 \rangle 5$. LET: $(\gamma_2, \dots, \gamma_k)$ and (n_2, \dots, n_k) be sequences such that $\gamma_k < \gamma_{k-1} < \cdots < \gamma_2$ and $\rho_1 = \omega^{\gamma_2} n_2 + \cdots + \omega^{\gamma_k} n_k$

$\langle 2 \rangle 6$. $\gamma_2 < \gamma_1$

PROOF: Since $\omega^{\gamma_2} \leq \rho_1 < \omega^{\gamma_1}$

$\langle 1 \rangle 2$. If

$$\begin{aligned} \gamma_k &< \gamma_{k-1} < \cdots < \gamma_1 \\ \gamma'_k &< \gamma'_{k-1} < \cdots < \gamma'_1 \end{aligned}$$

and

$$\omega^{\gamma_1} n_1 + \omega^{\gamma_2} n_2 + \cdots + \omega^{\gamma_k} n_k = \omega^{\gamma'_1} n'_1 + \omega^{\gamma'_2} n'_2 + \cdots + \omega^{\gamma'_k} n'_k$$

then $\gamma_i = \gamma'_i$ for all i and $n_i = n'_i$ for all i

PROOF: Prove by induction on i using the Logarithm Theorem.

□

Definition 8.8.2 (Cantor Normal Form). For any ordinal α , the *Cantor normal form* of α is the expression $\alpha = \omega^{\gamma_1} n_1 + \cdots + \omega^{\gamma_k} n_k$ such that n_1, \dots, n_k are nonzero natural numbers and $\gamma_k < \gamma_{k-1} < \cdots < \gamma_1$.

Chapter 9

The Cumulative Hierarchy

Definition 9.0.1 (Cumulative Hierarchy). Define the function $V : \mathbf{On} \rightarrow \mathbf{V}$ by transfinite recursion thus:

$$V_\alpha = \bigcup_{\beta < \alpha} \mathcal{P}V_\beta$$

Proposition 9.0.2. *For all $\alpha \in \mathbf{On}$, V_α is a transitive set.*

PROOF:

$\langle 1 \rangle 1.$ LET: $\alpha \in \mathbf{On}$

$\langle 1 \rangle 2.$ ASSUME: as transfinite induction hypothesis $\forall \beta < \alpha. V_\beta$ is a transitive set.

$\langle 1 \rangle 3.$ For all $\beta < \alpha$, $\mathcal{P}V_\beta$ is a transitive set.

PROOF: Proposition 1.6.4.

$\langle 1 \rangle 4.$ V_α is a transitive set.

PROOF: Proposition 1.6.3.

□

Proposition 9.0.3. *For any ordinals α and β , if $\beta < \alpha$ then $V_\beta \subseteq V_\alpha$.*

PROOF: Since $V_\beta \in \mathcal{P}V_\beta \subseteq V_\alpha$ and V_α is a transitive set. □

Theorem 9.0.4.

1. $V_0 = \emptyset$

2. $\forall \alpha \in \mathbf{On}. V_{\alpha+} = \mathcal{P}V_\alpha$

3. For any limit ordinal λ , $V_\lambda = \bigcup_{\alpha < \lambda} V_\alpha$.

PROOF:

$\langle 1 \rangle 1.$ $V_0 = \emptyset$

PROOF: Immediate from definition.

$\langle 1 \rangle 2.$ $\forall \alpha \in \mathbf{On}. V_{\alpha+} = \mathcal{P}V_\alpha$

PROOF:

- ⟨2⟩1. LET: $\alpha \in \mathbf{On}$
 ⟨2⟩2. For all $\beta < \alpha$ we have $\mathcal{P}V_\beta \subseteq \mathcal{P}V_\alpha$
 PROOF: Propositions 1.5.8 and 9.0.3.
 ⟨2⟩3. $V_{\alpha^+} = \mathcal{P}V_\alpha$

$$\begin{aligned}
 V_{\alpha^+} &= \bigcup_{\beta < \alpha^+} \mathcal{P}V_\beta \\
 &= \bigcup_{\beta < \alpha} \mathcal{P}V_\beta \cup \mathcal{P}V_\alpha \\
 &\quad \mathcal{P}V_\alpha
 \end{aligned}$$

□

- ⟨1⟩3. For any limit ordinal λ , $V_\lambda = \bigcup_{\alpha < \lambda} V_\alpha$

PROOF:

- ⟨2⟩1. $V_\lambda \subseteq \bigcup_{\alpha < \lambda} V_\alpha$

PROOF:

$$\begin{aligned}
 V_\lambda &= \bigcup_{\alpha < \lambda} \mathcal{P}V_\alpha \\
 &= \bigcup_{\alpha < \lambda} V_{\alpha^+} & (\langle 1 \rangle 2) \\
 &\subseteq \bigcup_{\alpha < \lambda} V_\alpha
 \end{aligned}$$

- ⟨2⟩2. $\bigcup_{\alpha < \lambda} V_\alpha \subseteq V_\lambda$

PROOF: Proposition 9.0.3.

□

Proposition 9.0.5. *For every set A , there exists an ordinal α such that $A \in V_\alpha$.*

PROOF:

- ⟨1⟩1. Let us say a set A is *grounded* iff there exists an ordinal α such that $A \in V_\alpha$.
 ⟨1⟩2. For any set A , if every element of A is grounded, then A is grounded.
 ⟨2⟩1. LET: A be a set.
 ⟨2⟩2. $S = \{\alpha \mid \exists a \in A. \alpha \text{ is the least ordinal such that } a \in V_\alpha\}$
 PROOF: S is a set by an Axiom of Replacement.
 ⟨2⟩3. LET: $\beta = \sup S$
 ⟨2⟩4. $A \subseteq V_\beta$
 ⟨3⟩1. LET: $a \in A$
 ⟨3⟩2. LET: α be the least ordinal such that $a \in V_\alpha$
 ⟨3⟩3. $\alpha \in S$
 ⟨3⟩4. $\alpha \leq \beta$
 ⟨3⟩5. $a \in V_\beta$
 ⟨2⟩5. $A \in V_{\beta^+}$
 ⟨1⟩3. ASSUME: for a contradiction there exists an ungrounded set.
 ⟨1⟩4. PICK a transitive set B that has an ungrounded member.
 PROOF: Pick a transitive set c , and take B to be the transitive closure of $\{c\}$.
 ⟨1⟩5. LET: $A = \{x \in B \mid x \text{ is ungrounded}\}$

- ⟨1⟩6. PICK $m \in A$ such that $m \cap A = \emptyset$
 PROOF: Axiom of Regularity.
 ⟨1⟩7. Every member of m is grounded.
 ⟨2⟩1. ASSUME: for a contradiction $x \in m$ is ungrounded.
 ⟨2⟩2. $x \in B$
 PROOF: Since B is transitive (⟨1⟩4).
 ⟨2⟩3. $x \in A$
 PROOF: ⟨1⟩5
 ⟨2⟩4. Q.E.D.
 PROOF: This contradicts ⟨1⟩6.
 ⟨1⟩8. m is grounded.
 PROOF: ⟨1⟩2
 ⟨1⟩9. Q.E.D.
 PROOF: This contradicts ⟨1⟩6.
 □

Definition 9.0.6 (Rank). The *rank* of a set A is the least ordinal α such that $A \in V_{\alpha+}$.

Proposition 9.0.7. For any set A we have

$$\text{rank } A = \bigcup_{a \in A} (\text{rank } a)^+$$

PROOF:

- ⟨1⟩1. LET: $\alpha = \bigcup_{a \in A} (\text{rank } a)^+$
 ⟨1⟩2. $A \subseteq V_\alpha$
 ⟨2⟩1. LET: $a \in A$
 ⟨2⟩2. $a \in V_{(\text{rank } a)^+}$
 ⟨2⟩3. $a \in V_\alpha$
 ⟨1⟩3. $A \in V_{\alpha+}$
 ⟨1⟩4. If $A \subseteq V_\beta$ then $\alpha \leq \beta$
 ⟨2⟩1. ASSUME: $A \subseteq V_\beta$
 ⟨2⟩2. For all $a \in A$ we have $(\text{rank } a)^+ \leq \beta$
 PROOF: Since $a \in V_\beta$.
 ⟨2⟩3. $\alpha \leq \beta$
 □

Corollary 9.0.7.1. For any sets a and b , if $a \in b$ then $\text{rank } a < \text{rank } b$.

Proposition 9.0.8. For any ordinal number α we have $\text{rank } \alpha = \alpha$.

PROOF:

- ⟨1⟩1. LET: α be an ordinal.
 ⟨1⟩2. ASSUME: as transfinite induction hypothesis $\forall \beta < \alpha. \text{rank } \beta = \beta$
 ⟨1⟩3. $\text{rank } \alpha = \bigcup_{\beta < \alpha} \beta^+$

PROOF:

$$\begin{aligned}\text{rank } \alpha &= \bigcup_{\beta < \alpha} (\text{rank } \beta)^+ \\ &= \bigcup_{\beta < \alpha} \beta^+\end{aligned}$$

$$\langle 1 \rangle 4. \bigcup_{\beta < \alpha} \beta^+ \leq \alpha$$

PROOF: Since for all $\beta < \alpha$ we have $\beta^+ \leq \alpha$.

$$\langle 1 \rangle 5. \alpha \leq \bigcup_{\beta < \alpha} \beta^+$$

$$\langle 2 \rangle 1. \text{ LET: } \gamma = \bigcup_{\beta < \alpha} \beta^+$$

$$\langle 2 \rangle 2. \text{ ASSUME: for a contradiction } \gamma < \alpha$$

$$\langle 2 \rangle 3. \gamma^+ \leq \bigcup_{\beta < \alpha} \beta^+ = \gamma$$

$$\langle 2 \rangle 4. \text{ Q.E.D.}$$

PROOF: This is a contradiction.

□

Definition 9.0.9 (Hereditarily Finite). A set is *hereditarily finite* iff it is in V_ω .

Chapter 10

Models of Set Theory

Definition 10.0.1 (Relativization). Let σ be a sentence in the language of set theory and \mathbf{M} a class. The *relativization* of σ to \mathbf{M} is the sentence $\sigma^{\mathbf{M}}$ formed by replacing every quantifier $\forall x$ with $\forall x \in \mathbf{M}$, and $\exists x$ with $\exists x \in \mathbf{M}$.

We write ' \mathbf{M} is a model of σ ' for the sentence $\sigma^{\mathbf{M}}$.

Theorem Schema 10.0.2. *For any class \mathbf{M} , the following is a theorem:*

If \mathbf{M} is a transitive class, then \mathbf{M} is a model of the Axiom of Extensionality.

PROOF:

$\langle 1 \rangle 1$. ASSUME: \mathbf{M} is a transitive class.

PROVE: $\forall x, y \in \mathbf{M} (\forall z \in \mathbf{M} (z \in x \Leftrightarrow z \in y) \Rightarrow x = y)$

$\langle 1 \rangle 2$. LET: $x, y \in \mathbf{M}$

$\langle 1 \rangle 3$. ASSUME: $\forall z \in \mathbf{M} (z \in x \Leftrightarrow z \in y)$

$\langle 1 \rangle 4$. $\forall z (z \in x \Leftrightarrow z \in y)$

PROOF: Since $z \in x \Rightarrow z \in \mathbf{M}$ and $z \in y \Rightarrow z \in \mathbf{M}$ by $\langle 1 \rangle 1$.

$\langle 1 \rangle 5$. $x = y$

□

Theorem 10.0.3. *If α is a non-zero ordinal then V_α is a model of the statement: The empty class is a set.*

PROOF:

$\langle 1 \rangle 1$. LET: $\alpha \neq 0$

PROVE: $\exists x \in V_\alpha. \forall y \in V_\alpha. y \notin x$

$\langle 1 \rangle 2$. $\emptyset \in V_\alpha$

$\langle 1 \rangle 3$. $\forall y \in V_\alpha. y \notin \emptyset$

□

Theorem 10.0.4. *For any limit ordinal λ , we have V_λ is a model of the statement: for any sets a and b , the class $\{a, b\}$ is a set.*

PROOF:

$\langle 1 \rangle 1$. LET: λ be a limit ordinal.

PROVE: $\forall a, b \in V_\lambda. \exists c \in V_\lambda. \forall x \in V_\lambda (x \in c \Leftrightarrow x = a \vee x = b)$
 (1)2. LET: $a, b \in V_\lambda$
 (1)3. PICK $\alpha, \beta < \lambda$ such that $a \in V_\alpha$ and $b \in V_\beta$
 (1)4. ASSUME: w.l.o.g. $\alpha \leq \beta$
 (1)5. $a, b \in V_\beta$
 (1)6. $\{a, b\} \in V_{\beta+1}$
 (1)7. $\{a, b\} \in V_\lambda$
 (1)8. $\forall x \in V_\lambda (x \in \{a, b\} \Leftrightarrow x = a \vee x = b)$
 \square

Theorem 10.0.5. *For any ordinal α , we have V_α is a model of the Union Axiom.*

PROOF:

(1)1. LET: α be an ordinal.
 PROVE: $\forall a \in V_\alpha. \exists b \in V_\alpha. \forall x \in V_\alpha (x \in b \Leftrightarrow \exists y \in V_\alpha (x \in y \wedge y \in a))$
 (1)2. LET: $a \in V_\alpha$
 (1)3. PICK $\beta < \alpha$ such that $a \subseteq V_\beta$
 (1)4. $\bigcup a \subseteq V_\beta$
 PROOF: V_β is a transitive set.
 (1)5. $\bigcup a \in V_\alpha$
 (1)6. $\forall x \in V_\alpha (x \in \bigcup a \Leftrightarrow \exists y \in V_\alpha (x \in y \wedge y \in a))$
 PROOF: V_α is a transitive set.
 \square

Theorem 10.0.6. *For any limit ordinal λ , we have V_λ is a model of the Power Set Axiom.*

PROOF:

(1)1. LET: λ be a limit ordinal.
 PROVE: $\forall a \in V_\lambda. \exists b \in V_\lambda. \forall x \in V_\lambda (x \in b \Leftrightarrow \forall y \in V_\lambda (y \in x \Rightarrow y \in a))$
 (1)2. LET: $a \in V_\lambda$
 (1)3. PICK $\alpha < \lambda$ such that $a \in V_\alpha$
 (1)4. $\mathcal{P}a \in V_{\alpha+1}$
 (1)5. $\mathcal{P}a \in V_\lambda$
 (1)6. $\forall x \in V_\lambda (x \in \mathcal{P}a \Leftrightarrow \forall y \in V_\lambda (y \in x \Rightarrow y \in a))$
 \square

Theorem Schema 10.0.7. *For any property $P[x, y_1, \dots, y_n]$, the following is a theorem:*

For any ordinal α , the set V_α is a model of the statement: for any sets a_1, \dots, a_n, B , the class $\{x \in B \mid P[x, a_1, \dots, a_n]\}$ is a set.

PROOF:

(1)1. LET: α be an ordinal.
 (1)2. LET: $a_1, \dots, a_n, B \in V_\alpha$
 (1)3. LET: $C = \{x \in B \mid P[x, a_1, \dots, a_n]^{V_\alpha}\}$
 (1)4. $C \in V_\alpha$

$\langle 1 \rangle 5. \forall x \in V_\alpha (x \in C \Leftrightarrow x \in B \wedge P[x, a_1, \dots, a_n]^{V_\alpha})$

□

Theorem 10.0.8. *For any ordinal $\alpha > \omega$, we have: V_α is a model of the Axiom of Infinity.*

PROOF:

$\langle 1 \rangle 1.$ LET: $\alpha > \omega$

$\langle 1 \rangle 2.$ $\mathbb{N} \in V_\alpha$

$\langle 1 \rangle 3.$ $\exists e \in V_\alpha (e \in \mathbb{N} \wedge \forall x \in V_\alpha. x \notin e)$

$\langle 1 \rangle 4.$ $\forall x \in V_\alpha (x \in \mathbb{N} \Rightarrow \exists y \in V_\alpha \forall z \in V_\alpha (z \in y \Leftrightarrow z \in x \vee z = x))$

□

Theorem 10.0.9. *For any ordinal α , we have V_α is a model of the Axiom of Choice.*

PROOF:

$\langle 1 \rangle 1.$ LET: α be an ordinal.

$\langle 1 \rangle 2.$ LET: $A \in V_\alpha$

$\langle 1 \rangle 3.$ ASSUME: $\forall x \in V_\alpha (x \in A \Rightarrow \exists y \in V_\alpha. y \in A)$

$\langle 1 \rangle 4.$ ASSUME: $\forall x, y, z \in V_\alpha (x \in A \wedge y \in A \wedge z \in x \wedge z \in y \Rightarrow x = y)$

$\langle 1 \rangle 5.$ A is a set of pairwise disjoint nonempty sets.

$\langle 1 \rangle 6.$ PICK c such that, for all $x \in A$, $x \cap c = \emptyset$

$\langle 1 \rangle 7.$ $c \cap \bigcup A \in V_\alpha$

$\langle 1 \rangle 8.$ $\forall x \in V_\alpha (x \in A \Rightarrow \exists y \in V_\alpha \forall z \in V_\alpha (z \in x \wedge z \in c \cap \bigcup A \Leftrightarrow z = y))$

□

Theorem 10.0.10. *For any ordinal α , we have V_α is a model of the Axiom of Regularity.*

PROOF:

$\langle 1 \rangle 1.$ LET: α be an ordinal.

$\langle 1 \rangle 2.$ LET: $A \in V_\alpha$

$\langle 1 \rangle 3.$ ASSUME: $\exists x \in V_\alpha. x \in A$

$\langle 1 \rangle 4.$ PICK $m \in A$ of least rank.

$\langle 1 \rangle 5.$ $m \in V_\alpha$

$\langle 1 \rangle 6.$ $\neg \exists x \in V_\alpha (x \in m \wedge x \in A)$

□

Theorem Schema 10.0.11. *For any axiom α of Zermelo set theory, the following is a theorem:*

For any limit ordinal $\lambda > \omega$, we have V_λ is a model of α .

PROOF: Theorems 10.0.2, 10.0.3, 10.0.4, 10.0.5, 10.0.6, 10.0.7, 10.0.8, 10.0.9, 10.0.10. □

Corollary Schema 10.0.11.1. *for any axiom α of Zermelo set theory, the following is a theorem:*

V_{ω_2} is a model of α .

Lemma 10.0.12. *There exists a well-ordered structure in V_{ω^2} whose ordinal is not in V_{ω^2} .*

PROOF: Take the well-ordered set $\mathbb{N} \times \{0, 1\}$ whose ordinal is ω^2 . \square

Corollary Schema 10.0.12.1. *There exists an instance α of the Axiom Schema of Replacement such that the following is a theorem:*

V_{ω^2} is not a model of α .

Chapter 11

Infinite Cardinals

11.1 Arithmetic of Infinite Cardinals

Proposition 11.1.1. *For any infinite cardinal κ we have $\kappa\kappa = \kappa$.*

PROOF:

- $\langle 1 \rangle 1$. PICK a set B with $|B| = \kappa$
- $\langle 1 \rangle 2$. LET: $\mathcal{H} = \{f \mid f = \emptyset \vee \exists A \subseteq B. (A \text{ is infinite} \wedge f : A \times A \approx A)\}$
- $\langle 1 \rangle 3$. For any chain $\mathcal{C} \subseteq \mathcal{H}$ we have $\bigcup \mathcal{C} \in \mathcal{H}$
 - $\langle 2 \rangle 1$. LET: $\mathcal{C} \subseteq \mathcal{H}$ be a chain.
 - $\langle 2 \rangle 2$. ASSUME: w.l.o.g. \mathcal{C} has a nonempty element.
 - $\langle 2 \rangle 3$. $\bigcup \mathcal{C}$ is a function.
 - $\langle 3 \rangle 1$. ASSUME: $(x, y), (x, z) \in \bigcup \mathcal{C}$
 - $\langle 3 \rangle 2$. PICK $f, g \in \mathcal{C}$ such that $f(x) = y$ and $g(x) = z$
 - $\langle 3 \rangle 3$. ASSUME: w.l.o.g. $f \subseteq g$
 - $\langle 3 \rangle 4$. $y = z$
 - $\langle 2 \rangle 4$. $\bigcup \mathcal{C}$ is injective.
- PROOF: Similar.
- $\langle 2 \rangle 5$. LET: $A = \text{ran } \bigcup \mathcal{C}$
- $\langle 2 \rangle 6$. A is infinite.
 - $\langle 3 \rangle 1$. PICK a nonzero $f \in \mathcal{C}$
 - $\langle 3 \rangle 2$. LET: A' be the infinite subset of B such that $f : A'^2 \approx A'$
 - $\langle 3 \rangle 3$. $A' \subseteq A$
- $\langle 2 \rangle 7$. $\text{dom } \bigcup \mathcal{C} = A^2$
 - $\langle 3 \rangle 1$. LET: $x, y \in A$
 - $\langle 3 \rangle 2$. PICK $f, g \in \mathcal{C}$ such that $x \in \text{ran } f$ and $y \in \text{ran } g$
 - $\langle 3 \rangle 3$. ASSUME: w.l.o.g. $f \subseteq g$
 - $\langle 3 \rangle 4$. LET: A' be the infinite subset of B such that $g : A'^2 \approx A'$
 - $\langle 3 \rangle 5$. $x, y \in A'$
 - $\langle 3 \rangle 6$. $(x, y) \in \text{dom } g$
 - $\langle 3 \rangle 7$. $(x, y) \in \text{dom } \bigcup \mathcal{C}$
- $\langle 2 \rangle 8$. $\bigcup \mathcal{C} \in \mathcal{H}$

- ⟨1⟩4. PICK a maximal $f_0 \in \mathcal{H}$
- ⟨1⟩5. $f_0 \neq \emptyset$
 - ⟨2⟩1. PICK a countably infinite subset A of B .
PROOF: Proposition 8.2.10.
 - ⟨2⟩2. PICK a bijection $f : A^2 \approx A$
PROOF: Proposition 8.2.9.
 - ⟨2⟩3. $\emptyset \subseteq f \in \mathcal{H}$
 - ⟨2⟩4. \emptyset is not maximal in \mathcal{H}
- ⟨1⟩6. LET: A_0 be the infinite subset of B such that $f_0 : A_0^2 \approx A_0$
- ⟨1⟩7. LET: $\lambda = |A_0|$
- ⟨1⟩8. λ is infinite.
- ⟨1⟩9. $\lambda^2 = \lambda$
- ⟨1⟩10. $\lambda = \kappa$
 - ⟨2⟩1. ASSUME: for a contradiction $\lambda < \kappa$
 - ⟨2⟩2. $\lambda \leq |B - A_0|$
 - ⟨2⟩3. PICK a subset $D \subseteq B - A_0$ with $|D| = \lambda$
 - ⟨2⟩4. $(A_0 \cup D)^2 = A_0^2 \cup (A_0 \times D) \cup (D \times A_0) \cup D^2$
 - ⟨2⟩5. LET: $C = (A_0 \times D) \cup (D \times A_0) \cup D^2$
 - ⟨2⟩6. $|C| = \lambda$
PROOF:

$$|(A_0 \times D) \cup (D \times A_0) \cup D^2| = \lambda^2 + \lambda^2 + \lambda^2$$

$$= \lambda + \lambda + \lambda \quad (\langle 1 \rangle 9)$$

$$= 3\lambda$$

$$\leq \lambda \cdot \lambda$$

$$= \lambda \quad (\langle 1 \rangle 9)$$
 - ⟨2⟩7. PICK a bijection $g : C \approx D$
 - ⟨2⟩8. $f_0 \cup g : (A_0 \cup D)^2 \approx A_0 \cup D$
 - ⟨2⟩9. Q.E.D.

PROOF: This contradicts the maximality of f_0 .

□

Theorem 11.1.2 (Absorption Law of Cardinal Arithmetic). *Let κ and λ be nonzero cardinal numbers such that at least one is infinite. Then*

$$\kappa + \lambda = \kappa\lambda = \max(\kappa, \lambda)$$

PROOF:

- ⟨1⟩1. ASSUME: w.l.o.g. $\lambda \leq \kappa$
- ⟨1⟩2. $\kappa + \lambda = \kappa\lambda = \kappa$

PROOF:

$$\begin{aligned}
 \kappa &\leq \kappa + \lambda \\
 &\leq \kappa + \kappa \\
 &= 2\kappa \\
 &\leq \kappa\lambda \\
 &\leq \kappa\kappa \\
 &= \kappa
 \end{aligned}$$

(Proposition 11.1.1)

□

11.2 Alephs

Definition 11.2.1 (Aleph). Let \aleph be the unique order isomorphism between **On** and the class of infinite cardinals.

Proposition 11.2.2. *The operation \aleph is normal.*

PROOF: Proposition 6.4.8 and Lemma 7.2.7. □

Definition 11.2.3 (Continuum Hypothesis). The *continuum hypothesis* is the statement that $\aleph_1 = 2^{\aleph_0}$.

Definition 11.2.4 (Generalised Continuum Hypothesis). The *generalised continuum hypothesis* is the statement that, for all α , $\aleph_{\alpha+} = 2^{\aleph_\alpha}$.

11.3 Beths

Definition 11.3.1 (Beth). Define the operation $\beth : \mathbf{On} \rightarrow \mathbf{Card}$ by transfinite recursion as follows:

$$\begin{aligned}
 \beth_0 &:= \aleph_0 \\
 \beth_{\alpha+} &:= 2^{\beth_\alpha} \\
 \beth_\lambda &:= \bigcup_{\alpha < \lambda} \beth_\alpha \quad (\lambda \text{ a limit ordinal})
 \end{aligned}$$

Proposition 11.3.2. *\beth is a normal operation.*

PROOF: It is continuous by definition, and $\beth_\alpha < \beth_{\alpha+}$ by Cantor's Theorem. □

Proposition 11.3.3. *The continuum hypothesis is equivalent to the statement $\beth_1 = \aleph_1$.*

The generalised continuum hypothesis is equivalent to the statement $\beth = \aleph$.

PROOF: Immediate from definitions. □

Lemma 11.3.4. *For any ordinal number α , we have $|V_{\omega+\alpha}| = \beth_\alpha$.*

PROOF:

(1)1. $|V_\omega| = \beth_0$

PROOF: Since V_ω is the union of \aleph_0 finite sets of increasing size.

(1)2. For any ordinal α , if $|V_{\omega+\alpha}| = \beth_\alpha$ then $|V_{\omega+\alpha+1}| = \beth_{\alpha+1}$

PROOF: Since $V_{\omega+\alpha+1} = \mathcal{P}V_{\omega+\alpha}$.

(1)3. For any limit ordinal λ , if $\forall \alpha < \lambda. |V_{\omega+\alpha}| = \beth_\alpha$ then $|V_{\omega+\lambda}| = \beth_\lambda$.

PROOF:

$$\begin{aligned} |V_{\omega+\lambda}| &= \left| \bigcup_{\alpha < \lambda} V_{\omega+\alpha} \right| \\ &= \sup_{\alpha < \lambda} |V_{\omega+\alpha}| \\ &= \sup_{\alpha < \lambda} \beth_\alpha \\ &= \beth_\lambda \end{aligned}$$

□

11.4 Cofinality

Definition 11.4.1 (Cofinal). Let λ be a limit ordinal and S a set of ordinals smaller than λ . Then S is *cofinal* in λ if and only if $\lambda = \sup S$.

Definition 11.4.2 (Cofinality). For any ordinal α , define the *cofinality* of α , $\text{cf } \alpha$, as follows:

- $\text{cf } 0 = 0$
- For any ordinal α , $\text{cf } \alpha^+ = 1$
- For any limit ordinal λ , $\text{cf } \lambda$ is the smallest cardinal such that there exists a set S of ordinals cofinal in λ with $|S| = \text{cf } \lambda$.

Definition 11.4.3 (Regular). A cardinal κ is *regular* iff $\text{cf } \kappa = \kappa$; otherwise it is *singular*.

Proposition 11.4.4. \aleph_0 is regular.

PROOF: \aleph_0 is not the supremum of $< \aleph_0$ smaller ordinals, because a finite union of finite ordinals is finite. □

Proposition 11.4.5. For every ordinal α , $\aleph_{\alpha+1}$ is regular.

PROOF: If S is a set of ordinals with $|S| < \aleph_{\alpha+1}$ and $\forall \beta \in S. \beta < \aleph_{\alpha+1}$, then we have $|S| \leq \aleph_\alpha$ and $\forall \beta \in S. \beta \leq \aleph_\alpha$, hence

$$\begin{aligned} \left| \bigcup S \right| &\leq \aleph_\alpha^2 && \text{(Proposition 7.2.6)} \\ &= \aleph_\alpha && \text{(Proposition 11.1.1)} \end{aligned}$$

Proposition Schema 11.4.6. For any class \mathbf{T} , the following is a theorem.

Assume $\mathbf{T} : \mathbf{On} \rightarrow \mathbf{On}$ is a normal operation. For any limit ordinal λ we have $\text{cf } \mathbf{T}(\lambda) = \text{cf } \lambda$.

PROOF:

- $\langle 1 \rangle 1.$ $\text{cf } \mathbf{T}(\lambda) \leq \text{cf } \lambda$
 - $\langle 2 \rangle 1.$ PICK a set S of ordinals $< \lambda$ with $|S| = \text{cf } \lambda$ and $\sup S = \lambda$
 - $\langle 2 \rangle 2.$ $\mathbf{T}(\lambda) = \sup_{\alpha \in S} \mathbf{T}(\alpha)$
PROOF: Theorem 6.4.7.
- $\langle 1 \rangle 2.$ $\text{cf } \lambda \leq \text{cf } \mathbf{T}(\lambda)$
 - $\langle 2 \rangle 1.$ PICK a set A of ordinals $< \mathbf{T}(\lambda)$ such that $|A| = \text{cf } \mathbf{T}(\lambda)$ and $\sup A = \mathbf{T}(\lambda)$
 - $\langle 2 \rangle 2.$ LET: $B = \{\gamma < \lambda \mid \exists \alpha \in A. |\alpha| = \mathbf{T}(\gamma)\}$
 - $\langle 2 \rangle 3.$ $|B| \leq |A| = \text{cf } \mathbf{T}(\lambda)$
PROVE: $\sup B = \lambda$
 - $\langle 2 \rangle 4.$ $\forall \alpha \in A. |\alpha| \leq \mathbf{T}(\sup B)$
 - $\langle 2 \rangle 5.$ $\forall \alpha \in A. \alpha < \mathbf{T}(\sup B + 1)$
 - $\langle 2 \rangle 6.$ $\aleph_\lambda = \sup A \leq \mathbf{T}(\sup B + 1)$
 - $\langle 2 \rangle 7.$ $\lambda \leq \sup B + 1$
 - $\langle 2 \rangle 8.$ $\lambda \leq \sup B$
PROOF: λ is a limit ordinal.
 - $\langle 2 \rangle 9.$ $\sup B = \lambda$

□

Corollary 11.4.6.1. \aleph_ω is singular.

PROOF: $\text{cf } \aleph_\omega = \text{cf } \aleph_0 = \aleph_0$. □

Corollary 11.4.6.2. The operation cf is not strictly monotone or continuous.

PROOF: $\text{cf } \aleph_\omega < \text{cf } \aleph_1$ □

Definition 11.4.7 (Weakly Inaccessible). A cardinal is *weakly inaccessible* iff it is \aleph_λ for some limit ordinal λ and regular.

Lemma 11.4.8. Let λ be a limit ordinal. Then there exists a strictly increasing $\text{cf } \lambda$ -sequence that converges to λ .

PROOF:

- $\langle 1 \rangle 1.$ PICK a set S of ordinals $< \lambda$ with $|S| = \text{cf } \lambda$ and $\sup S = \lambda$
- $\langle 1 \rangle 2.$ PICK a bijection $a : \text{cf } \lambda \approx S$
- $\langle 1 \rangle 3.$ PICK a strictly increasing subsequence $(b_\delta)_{\delta < \beta}$ of a that converges to λ .
PROOF: Lemma 6.6.5.
- $\langle 1 \rangle 4.$ $\beta = \text{cf } \lambda$
PROOF: By minimality of $\text{cf } \lambda$.

□

Corollary 11.4.8.1. Let λ be a limit ordinal. Then $\text{cf } \lambda$ is the least ordinal such that there exists a strictly increasing $\text{cf } \lambda$ -sequence that converges to λ .

Proposition 11.4.9. For any ordinal λ , $\text{cf } \lambda$ is a regular cardinal.

PROOF:

- (1)1. LET: λ be an ordinal.
 (1)2. ASSUME: w.l.o.g. λ is a limit ordinal.
 (1)3. PICK a strictly increasing sequence $(a_\alpha)_{\alpha < \text{cf } \lambda}$ that converges to λ .
 (1)4. LET: S be a set of ordinals $< \text{cf } \lambda$ such that $|S| = \text{cf } \text{cf } \lambda$ and $\sup S = \text{cf } \lambda$.
 (1)5. LET: $a(S) = \{a_\alpha \mid \alpha \in S\}$
 (1)6. $a(S)$ is cofinal in λ .
 (2)1. LET: $\beta < \lambda$
 (2)2. PICK $\gamma < \text{cf } \lambda$ such that $\beta < a_\gamma$
 (2)3. PICK $\delta \in S$ such that $\gamma < \delta$
 (2)4. $a_\delta \in a(S)$ and $\beta < a_\gamma < a_\delta$
 (1)7. $\text{cf } \lambda \leq \text{cf } \text{cf } \lambda$
 PROOF: Since $a(S)$ is a set of ordinals $< \lambda$ with $|a(S)| = \text{cf } \text{cf } \lambda$ and $\sup a(S) = \lambda$.
 (1)8. $\text{cf } \text{cf } \lambda = \text{cf } \lambda$
 □

Theorem 11.4.10. *Let λ be an infinite cardinal. Then $\text{cf } \lambda$ is the least cardinal such that λ can be partitioned into $\text{cf } \lambda$ sets, each of cardinality $< \lambda$.*

PROOF:

- (1)1. λ can be partitioned into $\text{cf } \lambda$ sets, each of cardinality $< \lambda$
 (2)1. PICK a strictly increasing sequence of ordinals $(a_\alpha)_{\alpha < \text{cf } \lambda}$ that converges to λ
 (2)2. $\{\{\beta \mid a_\alpha \leq \beta < a_{\alpha+1}\} \mid \alpha < \text{cf } \lambda\}$ is a partition of λ into $\text{cf } \lambda$ sets, each of cardinality $< \lambda$
 (1)2. If λ can be partitioned into κ sets, each of cardinality $< \lambda$, then $\text{cf } \lambda \leq \kappa$.
 (2)1. LET: \mathcal{A} be a partition of λ into sets of cardinality $< \lambda$
 (2)2. LET: $\kappa = |P|$
 (2)3. PICK a bijection $A : \kappa \approx P$
 (2)4. $\lambda = \bigcup_{\xi < \kappa} A(\xi)$
 (2)5. For all $\xi < \kappa$ we have $|A(\xi)| < \lambda$
 (2)6. LET: $\mu = \sup_{\xi < \kappa} |A(\xi)|$
 (2)7. $\mu \leq \lambda$
 (2)8. For all $\xi < \kappa$ we have $|A(\xi)| \leq \mu$
 (2)9. $\lambda \leq \mu\kappa$
 PROOF: Proposition 7.2.6.
 (2)10. ASSUME: w.l.o.g. $\kappa < \lambda$
 PROOF: If $\lambda \leq \kappa$ then $\text{cf } \lambda \leq \kappa$ since $\text{cf } \lambda \leq \lambda$.
 (2)11. $\lambda = \mu$
 PROOF:

$$\lambda \leq \mu\kappa \quad (\langle 2 \rangle 9)$$

$$\leq \lambda\lambda \quad (\langle 2 \rangle 7, \langle 2 \rangle 10)$$

$$= \lambda \quad (\text{Proposition 11.1.1})$$

- (2)12. $\{|A(\xi)| \mid \xi < \kappa\}$ is a set of $\leq \kappa$ ordinals all $< \lambda$ whose supremum is λ
 (2)13. $\text{cf } \lambda \leq \kappa$

□

Theorem 11.4.11 (König). *For any infinite cardinal κ we have $\kappa < \text{cf } 2^\kappa$.*

PROOF:

$\langle 1 \rangle 1$. ASSUME: for a contradiction $\text{cf } 2^\kappa \leq \kappa$

$\langle 1 \rangle 2$. LET: $S = 2^\kappa$

$\langle 1 \rangle 3$. PICK a partition $\{A_\xi \mid \xi < \kappa\}$ of S^κ with $\forall \xi < \kappa. |A_\xi| < 2^\kappa$.

PROOF: Theorem 11.4.10.

$\langle 1 \rangle 4$. $\forall \xi < \kappa. \{g(\xi) \mid g \in A_\xi\} \subsetneq S$

PROOF: We do not have equality because $|\{g(\xi) \mid g \in A_\xi\}| \leq |A_\xi| < 2^\kappa$.

$\langle 1 \rangle 5$. For all $\xi < \kappa$, choose $s_\xi \in S - \{g(\xi) \mid g \in A_\xi\}$

$\langle 1 \rangle 6$. $s \in S^\kappa$

$\langle 1 \rangle 7$. For all $\xi < \kappa$ we have $s \notin A_\xi$

PROOF: Since for all $\xi < \kappa$ and $g \in A_\xi$ we have $s_\xi(\xi) \neq g(\xi)$.

$\langle 1 \rangle 8$. Q.E.D.

PROOF: This contradicts $\langle 1 \rangle 3$.

□

Corollary 11.4.11.1.

$$2^{\aleph_0} \neq \aleph_\omega$$

Proposition 11.4.12. *For any ordinal α , we have $\text{cf } \alpha$ is the least cardinal such that α is the strict supremum of $\text{cf } \alpha$ smaller ordinals.*

PROOF:

$\langle 1 \rangle 1$. CASE: $\alpha = 0$

PROOF: Since $0 = \text{ssup } \emptyset$.

$\langle 1 \rangle 2$. CASE: $\alpha = \beta^+$

PROOF: Since $\beta^+ = \text{ssup}\{\beta\}$.

$\langle 1 \rangle 3$. CASE: α is a limit ordinal.

$\langle 2 \rangle 1$. There exists a set S of ordinals $< \alpha$ such that $|S| = \text{cf } \alpha$ and $\alpha = \text{ssup } S$.

$\langle 3 \rangle 1$. PICK a set S of ordinals $< \alpha$ such that $|S| = \text{cf } \alpha$ and $\sup S = \alpha$

PROVE: $\alpha = \text{ssup } S$

$\langle 3 \rangle 2$. $\forall \beta \in S. \beta < \alpha$

$\langle 3 \rangle 3$. For any ordinal γ , if $\forall \beta \in S. \beta < \gamma$ then $\alpha \leq \gamma$

$\langle 2 \rangle 2$. If T is a set of ordinals $< \alpha$ such that $\alpha = \text{ssup } T$, then $\text{cf } \alpha \leq |T|$.

$\langle 3 \rangle 1$. LET: T be a set of ordinals $< \alpha$ such that $\alpha = \text{ssup } T$

$\langle 3 \rangle 2$. $\alpha = \sup T$

$\langle 4 \rangle 1$. For all $\beta \in T$ we have $\beta \leq \alpha$

$\langle 4 \rangle 2$. LET: μ be any upper bound for T

PROVE: $\alpha \leq \mu$

$\langle 4 \rangle 3$. $\alpha \leq \mu + 1$

PROOF: Since $\forall \beta \in T. \beta < \mu + 1$.

$\langle 4 \rangle 4$. $\alpha \neq \mu + 1$

PROOF: Since α is a limit ordinal.

$\langle 4 \rangle 5$. $\alpha < \mu + 1$

$\langle 4 \rangle 6$. $\alpha \leq \mu$

$\langle 3 \rangle 3$. $\text{cf } \alpha \leq |T|$

□

11.5 Inaccessible Cardinals

Definition 11.5.1 (Inaccessible Cardinal). A cardinal number κ is *inaccessible* iff:

- $\kappa > \aleph_0$
- $\forall \lambda < \kappa. 2^\lambda < \kappa$ (cardinal exponentiation)
- κ is regular.

Any inaccessible cardinal is weakly inaccessible.

PROOF:

$\langle 1 \rangle 1$. LET: $\kappa = \aleph_\lambda$ be weakly inaccessible.

PROVE: λ is a limit ordinal.

$\langle 1 \rangle 2$. $\lambda \neq 0$

$\langle 1 \rangle 3$. ASSUME: for a contradiction $\lambda = \beta + 1$

$\langle 1 \rangle 4$. $\aleph_\beta < \kappa$

$\langle 1 \rangle 5$. $2^{\aleph_\beta} < \kappa$

$\langle 1 \rangle 6$. $\aleph_{\beta+1} < \kappa$

PROOF: Since $\aleph_{\beta+1} \leq 2^{\aleph_\beta}$.

$\langle 1 \rangle 7$. Q.E.D.

PROOF: This contradicts $\langle 1 \rangle 3$.

□

Proposition 11.5.2. *If the Generalized Continuum Hypothesis is true, then every weakly inaccessible cardinal is inaccessible.*

PROOF:

$\langle 1 \rangle 1$. ASSUME: The Generalized Continuum Hypothesis.

$\langle 1 \rangle 2$. LET: $\kappa = \aleph_\lambda$ be weakly inaccessible.

$\langle 1 \rangle 3$. $\kappa > \aleph_0$

PROOF: $\lambda > 0$ because λ is a limit ordinal.

$\langle 1 \rangle 4$. For all $\mu < \kappa$ we have $2^\mu < \kappa$

$\langle 2 \rangle 1$. LET: $\mu < \kappa$

$\langle 2 \rangle 2$. LET: $\mu = \aleph_\alpha$

$\langle 2 \rangle 3$. $\alpha < \lambda$

$\langle 2 \rangle 4$. $\alpha + 1 < \lambda$

PROOF: λ is a limit ordinal.

$\langle 2 \rangle 5$. $2^\mu < \kappa$

PROOF:

$$\begin{aligned}
 2^\mu &= 2^{\aleph_\alpha} && (\langle 2 \rangle 2) \\
 &= 2^{\beth_\alpha} && (\langle 1 \rangle 1) \\
 &= \beth_{\alpha+1} \\
 &= \aleph_{\alpha+1} && (\langle 1 \rangle 1) \\
 &< \aleph_\lambda && (\langle 2 \rangle 4) \\
 &= \kappa && (\langle 1 \rangle 2)
 \end{aligned}$$

$\langle 1 \rangle 5.$ κ is regular.

PROOF: $\langle 1 \rangle 2$

□

Lemma 11.5.3. *Let κ be an inaccessible cardinal. For every ordinal $\alpha < \kappa$ we have $\beth_\alpha < \kappa$.*

PROOF:

$\langle 1 \rangle 1.$ $\beth_0 < \kappa$

PROOF: Since $\kappa > \aleph_0$.

$\langle 1 \rangle 2.$ For any ordinal α , if $\beth_\alpha < \kappa$ then $\beth_{\alpha+1} < \kappa$.

PROOF: Since $\beth_{\alpha+1} = 2^{\beth_\alpha} < \kappa$.

$\langle 1 \rangle 3.$ For any limit ordinal λ , if $\forall \alpha < \lambda. \beth_\alpha < \kappa$ and $\lambda < \kappa$ then $\beth_\lambda < \kappa$.

PROOF: By regularity of κ , since \beth_λ is the union of $|\lambda|$ cardinals all $< \kappa$.

□

Lemma 11.5.4. *Let κ be an inaccessible cardinal. For all $A \in V_\kappa$ we have $|A| < \kappa$.*

PROOF:

$\langle 1 \rangle 1.$ LET: $A \in V_\kappa$

$\langle 1 \rangle 2.$ PICK $\alpha < \kappa$ such that $A \in V_\alpha$

$\langle 1 \rangle 3.$ $A \subseteq V_\alpha$

$\langle 1 \rangle 4.$ $|A| \leq |V_\alpha| \leq \beth_\alpha < \kappa$

□

Theorem Schema 11.5.5. *For every axiom α of ZFC, the following is a theorem:*

For any inaccessible cardinal κ , we have V_κ is a model of α .

PROOF: For every axiom except the Replacement Axioms, we have Corollary 10.0.11.1.

For an Axiom of Replacement using the property $P[x, y, z_1, \dots, z_n]$, we reason as follows:

$\langle 1 \rangle 1.$ LET: κ be an inaccessible cardinal

PROVE:

$$\begin{aligned} & \forall a_1, \dots, a_n, B \in V_\kappa (\forall x \in B. \forall y, y' \in V_\kappa \\ & (P[x, y, a_1, \dots, a_n]^{V_\kappa} \wedge P[x, y', a_1, \dots, a_n]^{V_\kappa} \Rightarrow y = y') \Rightarrow \\ & \exists C \in V_\kappa \forall y \in V_\kappa (y \in C \Leftrightarrow \exists x \in B. P[x, y, a_1, \dots, a_n]^{V_\kappa})) \end{aligned}$$

$\langle 1 \rangle 2.$ LET: $a_1, \dots, a_n, B \in V_\kappa$

$\langle 1 \rangle 3.$ ASSUME: for all $x \in B$, there exists at most one $y \in V_\kappa$ such that

$$P[x, y, a_1, \dots, a_n]^{V_\kappa}.$$

$\langle 1 \rangle 4.$ LET: $F = \{(x, y) \in B \times V_\kappa \mid P[x, y, a_1, \dots, a_n]^{V_\kappa}\}$

$\langle 1 \rangle 5.$ LET: $C = \text{ran } F$

PROVE: $C \in V_\kappa$

$\langle 1 \rangle 6.$ LET: $S = \{\text{rank } F(x) \mid x \in \text{dom } F\}$

$\langle 1 \rangle 7.$ $|S| < \kappa$

PROOF: Since $|S| \leq |\text{dom } F| \leq |B| < \kappa$.

$\langle 1 \rangle 8. \forall \alpha \in S. \alpha < \kappa$

PROOF: Since $F(x) \in V_\kappa$ for all $x \in \text{dom } F$.

$\langle 1 \rangle 9. \sup S < \kappa$

PROOF: Since κ is regular.

$\langle 1 \rangle 10. \text{rank } C \leq \sup S + 1$

$\langle 1 \rangle 11. \text{rank } C < \kappa$

$\langle 1 \rangle 12. C \in V_\kappa$

□

Chapter 12

Group Theory

12.1 Groups

Definition 12.1.1 (Group). A *group* G consists of a set G and a function $\cdot : G^2 \rightarrow G$ such that:

1. \cdot is associative
2. There exists $e \in G$ such that $\forall x \in G. xe = x$ and $\forall x \in G. \exists y \in G. xy = e$.

Proposition 12.1.2. *The inverse of an element in a group is unique.*

PROOF:

$\langle 1 \rangle$ 1. ASSUME: b and b' are inverses of a .

$\langle 1 \rangle$ 2. $b = b'$

PROOF:

$$\begin{aligned} b &= be \\ &= bab' \\ &= eb' \\ &= b' \end{aligned}$$

□

Definition 12.1.3. We write x^{-1} for the inverse of x .

Proposition 12.1.4. *In any group, if $ab = ac$ then $b = c$.*

PROOF:

$$\begin{aligned} b &= eb \\ &= a^{-1}ab \\ &= a^{-1}ac \\ &= ec \\ &= c \end{aligned}$$

□

12.2 Abelian Groups

Definition 12.2.1 (Abelian group). An *Abelian group* is a group whose multiplication is commutative.

We may say we are writing an Abelian group *additively*, meaning we write $a + b$ for ab , 0 for e and $-a$ for a^{-1} . In this case we write $a - b$ for ab^{-1} .

Chapter 13

Ring Theory

13.1 Rings

Definition 13.1.1 (Commutative Ring). A *commutative ring* consists of a set R and two binary operations $+$, \cdot on R such that:

- D is an Abelian group under $+$. Let us write 0 for its identity element.
- \cdot is commutative and associative, and distributes over $+$.
- \cdot has an identity element 1 that is different from 0 .

Proposition 13.1.2. *In any commutative ring, $0x = 0$.*

PROOF:

$$\begin{aligned}(0 + 0)x &= 0x \\ \therefore 0x + 0x &= 0x + 0 \\ \therefore 0x &= 0 && \text{(Proposition 12.1.4)} \square\end{aligned}$$

Proposition 13.1.3. *In any commutative ring, $(-a)b = -(ab)$.*

PROOF:

$$\begin{aligned}ab + (-a)b &= (a + (-a))b \\ &= 0b \\ &= 0 && \text{(Proposition 13.1.2)} \square\end{aligned}$$

13.2 Ordered Rings

Definition 13.2.1 (Ordered Commutative Ring). An *ordered commutative ring* consists of a commutative ring R with a linear order $<$ on R such that:

- for all $x, y, z \in R$, we have $x < y$ if and only if $x + z < y + z$.

- for all $x, y, z \in R$, if $0 < z$ then we have $x < y$ if and only if $xz < yz$.

Proposition 13.2.2. *In any ordered commutative ring, $0 < 1$.*

PROOF: If $1 < 0$ then we have $0 < -1$ and so $0 < (-1)(-1) = 1$, which is a contradiction. \square

Proposition 13.2.3. *The ordering on an ordered commutative ring is dense; that is, if $x < y$ then there exists z such that $x < z < y$.*

PROOF: Take $z = (x + y)/2$. \square

13.3 Integral Domains

Definition 13.3.1 (Integral Domain). An *integral domain* is a commutative ring such that, for all $a, b \in D$, if $ab = 0$ then $a = 0$ or $b = 0$.

Proposition 13.3.2. *In any integral domain, if $ab = ac$ and $a \neq 0$ then $b = c$.*

PROOF: We have $a(b - c) = 0$ and $a \neq 0$ so $b - c = 0$ hence $b = c$. \square

Definition 13.3.3 (Ordered Integral Domain). An *ordered integral domain* is an ordered commutative ring that is an integral domain.

Chapter 14

Field Theory

14.1 Fields

Definition 14.1.1 (Field). A *field* F is a commutative ring such that $0 \neq 1$ and, for all $x \in F$, if $x \neq 0$ then there exists $y \in F$ such that $xy = 1$.

Proposition 14.1.2. *Every field is an integral domain.*

PROOF: If $ab = 0$ and $a \neq 0$ then $b = a^{-1}ab = 0$. \square

Proposition 14.1.3. *In any field F , we have $F - \{0\}$ is an Abelian group under multiplication.*

PROOF: Immediate from the definition. \square

Definition 14.1.4 (Field of Fractions). Let D be an integral domain. The *field of fractions* of D is the quotient set $F = (D \times (D - \{0\})) / \sim$ where

$$(a, b) \sim (c, d) \Leftrightarrow ad = bc$$

under

$$\begin{aligned} [(a, b)] + [(c, d)] &= [(ad + bc, bd)] \\ [(a, b)][(c, d)] &= [(ac, bd)] \end{aligned}$$

We prove this is a field.

PROOF:

$\langle 1 \rangle 1.$ \sim is an equivalence relation on $D \times (D - \{0\})$.

PROOF:

$\langle 2 \rangle 1.$ \sim is reflexive.

PROOF: We always have $ab = ba$.

$\langle 2 \rangle 2.$ \sim is symmetric.

PROOF: If $ad = bc$ then $cb = da$.

$\langle 2 \rangle 3$. \sim is transitive.

$\langle 3 \rangle 1$. ASSUME: $(a, b) \sim (c, d) \sim (e, f)$

$\langle 3 \rangle 2$. $ad = bc$ and $cf = de$

$\langle 3 \rangle 3$. $adf = bde$

PROOF: $adf = bcf = bde$

$\langle 3 \rangle 4$. $af = be$

PROOF: Proposition 13.3.2.

□

$\langle 1 \rangle 2$. Addition is well-defined.

PROOF:

$\langle 2 \rangle 1$. If $b \neq 0$ and $d \neq 0$ then $bd \neq 0$.

PROOF: Since D is an integral domain.

$\langle 2 \rangle 2$. If $ab' = a'b$ and $cd' = c'd$ then $(ad + bc)b'd' = (a'd' + b'c')bd$.

PROOF:

$$\begin{aligned} (ad + bc)b'd' &= ab'dd' + bb'cd' \\ &= a'bdd' + bb'c'd \\ &= (a'd' + b'c')bd \end{aligned}$$

□

$\langle 1 \rangle 3$. Multiplication is well-defined.

PROOF:

$\langle 2 \rangle 1$. If $b \neq 0$ and $d \neq 0$ then $bd \neq 0$.

PROOF: Since D is an integral domain.

$\langle 2 \rangle 2$. If $[(a, b)] = [(a', b')]$ and $[(c, d)] = [(c', d')]$ then $[(ac, bd)] = [(a'c', b'd')]$.

PROOF: If $ab' = a'b$ and $cd' = c'd$ then $acb'd' = a'c'bd$.

□

$\langle 1 \rangle 4$. Addition is commutative.

PROOF: $[(a, b)] + [(c, d)] = [(ad + bc, bd)] = [(cb + da, db)] = [(c, d)] + [(a, b)]$ □

$\langle 1 \rangle 5$. Addition is associative.

PROOF:

$$\begin{aligned} [(a, b)] + ([[(c, d)] + [(e, f)]] &= [(a, b)] + [(cf + de, df)] \\ &= [(adf + bcf + bde, bdf)] \\ &= [(ad + bc, bd)] + [(e, f)] \\ &= ([[(a, b)] + [(c, d)]] + [(e, f)]) \quad \square \end{aligned}$$

$\langle 1 \rangle 6$. For any $x \in F$ we have $x + [(0, 1)] = x$

PROOF: $[(a, b)] + [(0, 1)] = [(a \cdot 1 + b \cdot 0, b \cdot 1)] = [(a, b)]$ □

$\langle 1 \rangle 7$. For any $x \in F$, there exists $y \in F$ such that $x + y = [(0, 1)]$.

PROOF: $[(a, b)] + [(-a, b)] = [(ab - ab, b^2)] = [(0, b^2)] = [(0, 1)]$ □

$\langle 1 \rangle 8$. Multiplication is commutative.

PROOF: $[(a, b)][(c, d)] = [(c, d)][(a, b)] = [(ac, bd)]$. □

$\langle 1 \rangle 9$. Multiplication is associative.

PROOF: $[(a, b)]([[(c, d)][(e, f)]] = ([[(a, b)][(c, d)]][(e, f)] = [(ace, bdf)]$. □

$\langle 1 \rangle 10$. For any $x \in F$ we have $x[(1, 1)] = x$

PROOF: $[(a, b)][(1, 1)] = [(a, b)]$ □

$\langle 1 \rangle 11$. For any non-zero $x \in F$, there exists $y \in F$ such that $xy = [(1, 1)]$.

PROOF:

- $\langle 2 \rangle 1$. LET: $[(a, b)] \in \mathbb{Q}$
- $\langle 2 \rangle 2$. ASSUME: $[(a, b)] \neq [(0, 1)]$
- $\langle 2 \rangle 3$. $a \neq 0$
- $\langle 2 \rangle 4$. $[(a, b)][(b, a)] = [(1, 1)]$

□

□

Definition 14.1.5. For any field F , let $N(F)$ be the intersection of all the subsets $S \subseteq F$ such that $1 \in S$ and $\forall x \in S. x + 1 \in S$.

Definition 14.1.6 (Characteristic Zero). A field F has *characteristic 0* iff $0 \notin N(F)$.

Proposition 14.1.7. In a field F with characteristic 0, the function $n : \mathbb{N} \rightarrow N(F)$ defined by

$$\begin{aligned} n(0) &= 1 \\ n(x + 1) &= n(x) + 1 \end{aligned}$$

is a bijection.

PROOF:

- $\langle 1 \rangle 1$. n is injective.
- $\langle 2 \rangle 1$. ASSUME: for a contradiction $n(i) = n(j)$ with $i \neq j$
- $\langle 2 \rangle 2$. ASSUME: w.l.o.g. $i < j$
- $\langle 2 \rangle 3$. $n(j - i) = 0$
- $\langle 2 \rangle 4$. Q.E.D.

PROOF: This contradicts the fact that F has characteristic 0.

- $\langle 1 \rangle 2$. n is surjective.

PROOF: Since $\text{ran } n$ is a subset of F that includes 1 and is closed under $+1$.

□

Definition 14.1.8. In any field F , let

$$I(F) = N(F) \cup \{0\} \cup \{-x \mid x \in N(F)\}$$

Definition 14.1.9. In any field F , let

$$Q(F) = \{x/y \mid x, y \in I(F), y \neq 0\}$$

Proposition 14.1.10. $Q(F)$ is the smallest subfield of F .

PROOF: $Q(F)$ is closed under $+$ and \cdot , and any subset of F closed under $+$ and \cdot that contains 0 and 1 must include $Q(F)$. □

Theorem 14.1.11. Let F and G be fields of characteristic 0. Then there exists a unique field isomorphism between $Q(F)$ and $Q(G)$.

PROOF:

- (1)1. LET: $\phi : N(F) \rightarrow N(G)$ be the unique function such that $\phi(1) = 1$ and $\forall x \in N(F). \phi(x+1) = \phi(x) + 1$.
- (1)2. ϕ is a bijection.
 PROOF: Similar to Proposition 14.1.7.
- (1)3. $\forall x, y \in N(F). \phi(x+y) = \phi(x) + \phi(y)$
 PROOF: Induction on y .
- (1)4. $\forall x, y \in N(F). \phi(xy) = \phi(x)\phi(y)$
 PROOF: Induction on y .
- (1)5. Extend ϕ to a bijection $I(F) \cong I(G)$ such that $\forall x, y \in I(F). \phi(x+y) = \phi(x) + \phi(y)$ and $\forall x, y \in I(F). \phi(xy) = \phi(x)\phi(y)$
- (2)1. Define $\phi(0) = 0$ and $\phi(-x) = -\phi(x)$ for $x \in N(F)$
- (3)1. $0 \notin N(F)$
- (3)2. For all $x \in N(F)$ we have $-x \notin N(F)$
 PROOF: Then we would have $x + -x = 0 \in N(F)$.
- (3)3. For all $x \in N(F)$ we have $-x \neq 0$
- (2)2. For all $x, y \in I(F)$ we have $\phi(x+y) = \phi(x) + \phi(y)$
 PROOF: Case analysis on x and y .
- (2)3. For all $x, y \in I(F)$ we have $\phi(xy) = \phi(x)\phi(y)$
 PROOF: Case analysis on x and y .
- (1)6. Extend ϕ to a bijection $Q(F) \cong Q(G)$ such that $\forall x, y \in Q(F). \phi(x+y) = \phi(x) + \phi(y)$ and $\forall x, y \in Q(F). \phi(xy) = \phi(x)\phi(y)$
- (2)1. Define $\phi(x/y) = \phi(x)/\phi(y)$
- (1)7. ϕ is unique.
- (2)1. LET: θ satisfy the theorem.
- (2)2. For all $x \in N(F)$ we have $\theta(x) = \phi(x)$
- (2)3. For all $x \in I(F)$ we have $\theta(x) = \phi(x)$
- (2)4. For all $x \in Q(F)$ we have $\theta(x) = \phi(x)$

□

14.2 Ordered Fields

Definition 14.2.1 (Ordered Field). An *ordered field* is an ordered commutative ring that is a field.

Proposition 14.2.2. Every ordered field F has characteristic 0.

PROOF: We have $0 < n$ for all $n \in N(F)$. □

Proposition 14.2.3. Let F be a field of characteristic 0. Then there exists a unique relation $<$ on $Q(F)$ that makes $Q(F)$ into an ordered field.

PROOF: Easy. □

Corollary 14.2.3.1. Let F and G be ordered fields. Let ϕ be the unique field isomorphism between $Q(F)$ and $Q(G)$. Then ϕ is an ordered field isomorphism.

Definition 14.2.4 (Archimedean). An ordered field F is *Archimedean* iff

$$\forall x \in F. \exists n \in N(F). n > x .$$

Proposition 14.2.5. *Let F be an Archimedean ordered field. Let $x, y \in F$ with $x > 0$. Then there exists $n \in N(F)$ such that $nx > y$.*

PROOF: Pick $n > y/x$. \square

Proposition 14.2.6. *Let F be an Archimedean ordered field. For all $x, y \in F$, if $x < y$, then there exists $r \in Q(F)$ such that $x < r < y$.*

PROOF:

$\langle 1 \rangle 1$. CASE: $x > 0$

$\langle 2 \rangle 1$. PICK $n \in N(F)$ such that $n(y - x) > 1$

PROOF: Proposition 14.2.5.

$\langle 2 \rangle 2$. $ny > 1 + nx$

$\langle 2 \rangle 3$. LET: m be the least element of $N(F)$ such that $m > nx$.

$\langle 2 \rangle 4$. $m - 1 \leq nx$

$\langle 2 \rangle 5$. $nx < m < ny$

$\langle 2 \rangle 6$. $x < m/n < y$

$\langle 1 \rangle 2$. CASE: $x \leq 0$

$\langle 2 \rangle 1$. PICK $k \in N(F)$ such that $k > -x$

$\langle 2 \rangle 2$. $0 < x + k < y + k$

$\langle 2 \rangle 3$. PICK $r \in Q(F)$ such that $x + k < r < y + k$

PROOF: $\langle 1 \rangle 1$

$\langle 2 \rangle 4$. $x < r - k < y$

Definition 14.2.7 (Complete). An ordered field F is *complete* iff every nonempty subset of F bounded above has a least upper bound.

Proposition 14.2.8. *Every complete ordered field is Archimedean.*

PROOF:

$\langle 1 \rangle 1$. LET: F be a complete ordered field.

$\langle 1 \rangle 2$. LET: $x \in F$

$\langle 1 \rangle 3$. ASSUME: for a contradiction there is no member of $N(F)$ greater than x .

$\langle 1 \rangle 4$. x is an upper bound for $N(F)$.

$\langle 1 \rangle 5$. LET: $y = \sup N(F)$

$\langle 1 \rangle 6$. PICK $n \in N(F)$ such that $y - 1 < n$

$\langle 1 \rangle 7$. $y < n + 1$

$\langle 1 \rangle 8$. Q.E.D.

PROOF: This is a contradiction.

\square

Proposition 14.2.9. *Let F be a complete ordered field and $a \in F$ be nonnegative. Then there exists $b \in F$ such that $b^2 = a$.*

PROOF:

$\langle 1 \rangle 1$. LET: $B = \{x \in F \mid 0 \leq x \leq 1 + a\}$

$\langle 1 \rangle 2$. LET: $\phi : B \rightarrow B$ be the function

$$\phi(x) = x + \frac{1}{2(1+a)}(a - x^2) .$$

- ⟨1⟩3. ϕ is strictly monotone.
 ⟨2⟩1. LET: $0 \leq x < y \leq 1 + a$
 ⟨2⟩2. $1 - \frac{x+y}{2(1+a)} > 0$
 ⟨2⟩3. $\phi(y) - \phi(x) = (y - x)(1 - \frac{x+y}{2(1+a)}) > 0$
 ⟨2⟩4. $\phi(x) < \phi(y)$
 ⟨1⟩4. PICK $b \in B$ such that $\phi(b) = b$.
 PROOF: Knaster Fixed-Point Theorem.
 ⟨1⟩5. $b^2 = a$
 \square

Theorem 14.2.10 (Uniqueness of the Complete Ordered Field). *If F and G are complete ordered fields, then there exists a unique bijection $\phi : F \cong G$ such that, for all $x, y \in F$,*

$$\begin{aligned}\phi(x + y) &= \phi(x) + \phi(y) \\ \phi(xy) &= \phi(x)\phi(y)\end{aligned}$$

This bijection also satisfies: for all $x, y \in F$,

$$x < y \Leftrightarrow \phi(x) < \phi(y) .$$

PROOF:

- ⟨1⟩1. PICK a bijection $\phi : Q(F) \cong Q(G)$ such that, for all $x, y \in Q(F)$,
 $\phi(x + y) = \phi(x) + \phi(y)$
 $\phi(xy) = \phi(x)\phi(y)$
 $x < y \Leftrightarrow \phi(x) < \phi(y)$

PROOF: Corollary 14.2.3.1.

- ⟨1⟩2. $Q(F)$ intersects every interval in F .

PROOF: Proposition 14.2.6.

- ⟨1⟩3. $Q(G)$ intersects every interval in G .

PROOF: Proposition 14.2.6.

- ⟨1⟩4. PICK an order isomorphism $\psi : F \cong G$ that extends ϕ .

PROOF: Theorem 5.1.21.

- ⟨1⟩5. $\forall x, y \in F. \psi(x + y) = \psi(x) + \psi(y)$
 ⟨2⟩1. LET: $x, y \in F$
 ⟨2⟩2. $\psi(x) + \psi(y) \not\leq \psi(x + y)$
 ⟨3⟩1. ASSUME: for a contradiction $\psi(x) + \psi(y) < \psi(x + y)$
 ⟨3⟩2. PICK $r' \in Q(G)$ such that $\psi(x) < r' < \psi(x + y) - \psi(y)$
 ⟨3⟩3. PICK $s' \in Q(G)$ such that $\psi(y) < s' < \psi(x + y) - r'$
 ⟨3⟩4. $r' + s' < \psi(x + y)$
 ⟨3⟩5. PICK $r, s \in Q(F)$ such that $\phi(r) = r'$ and $\phi(s) = s'$
 ⟨3⟩6. $\phi(r + s) = r' + s'$
 ⟨3⟩7. $\psi(x) < \psi(r)$
 ⟨3⟩8. $\psi(y) < \psi(s)$
 ⟨3⟩9. $\psi(x + y) > \psi(r + s)$
 ⟨3⟩10. $x < r$

- $\langle 3 \rangle 11. y < s$
- $\langle 3 \rangle 12. x + y > r + s$
- $\langle 3 \rangle 13. \text{Q.E.D.}$

PROOF: This is a contradiction.

- $\langle 2 \rangle 3. \psi(x + y) \not\leq \psi(x) + \psi(y)$

PROOF: Similar.

- $\langle 1 \rangle 6. \forall x, y \in F. \psi(xy) = \psi(x)\psi(y)$

- $\langle 2 \rangle 1. \text{LET: } x, y \in F$

- $\langle 2 \rangle 2. \text{CASE: } x \text{ and } y \text{ are positive.}$

- $\langle 3 \rangle 1. \psi(x)\psi(y) \not\leq \psi(xy)$

- $\langle 4 \rangle 1. \text{ASSUME: for a contradiction } \psi(x)\psi(y) < \psi(xy)$

- $\langle 4 \rangle 2. \text{PICK } r' \in Q(G) \text{ such that } \psi(x) < r' < \psi(xy)/\psi(y)$

- $\langle 4 \rangle 3. \text{PICK } s' \in Q(G) \text{ such that } \psi(y) < s' < \psi(xy)/r'$

- $\langle 4 \rangle 4. r's' < \psi(xy)$

- $\langle 4 \rangle 5. \text{PICK } r, s \in Q(F) \text{ such that } \phi(r) = r' \text{ and } \phi(s) = s'$

- $\langle 4 \rangle 6. \phi(rs) = r's'$

- $\langle 4 \rangle 7. x < r, y < s \text{ and } rs < xy$

- $\langle 4 \rangle 8. \text{Q.E.D.}$

PROOF: This is a contradiction.

- $\langle 3 \rangle 2. \psi(xy) \not\leq \psi(x)\psi(y)$

PROOF: Similar.

- $\langle 2 \rangle 3. \text{CASE: } x \text{ and } y \text{ are not both positive.}$

PROOF: Follows from $\langle 2 \rangle 2$ since $\psi(-x) = -\psi(x)$ by $\langle 1 \rangle 5$.

- $\langle 1 \rangle 7. \text{For any field isomorphism } \theta : F \cong G, \text{ we have } \theta = \psi.$

- $\langle 2 \rangle 1. \theta \upharpoonright Q(F) = \phi$

PROOF: Theorem 14.1.11.

- $\langle 2 \rangle 2. \theta \text{ is strictly monotone.}$

- $\langle 3 \rangle 1. \text{LET: } x, y \in F \text{ with } x < y$

- $\langle 3 \rangle 2. y - x > 0$

- $\langle 3 \rangle 3. \text{PICK } z \in F \text{ such that } z^2 = y - x$

- $\langle 3 \rangle 4. \theta(z)^2 = \theta(y) - \theta(x)$

- $\langle 3 \rangle 5. \theta(y) - \theta(x) > 0$

- $\langle 3 \rangle 6. \theta(x) < \theta(y)$

- $\langle 2 \rangle 3. \theta = \psi$

PROOF: By the uniqueness of ψ .

□

Chapter 15

Number Systems

15.1 The Integers

Definition 15.1.1. The set of *integers* \mathbb{Z} is the quotient set \mathbb{N}^2 / \sim , where $(m, n) \sim (p, q)$ iff $m + q = n + p$.

We prove \sim is an equivalence relation on \mathbb{N}^2 .

PROOF:

$\langle 1 \rangle 1.$ \sim is reflexive.

PROOF: For all $m, n \in \mathbb{N}$ we have $m + n = n + m$.

$\langle 1 \rangle 2.$ \sim is symmetric.

PROOF: If $m + q = n + p$ then $p + n = q + m$.

$\langle 1 \rangle 3.$ \sim is transitive.

$\langle 2 \rangle 1.$ ASSUME: $(m, n) \sim (p, q) \sim (r, s)$

$\langle 2 \rangle 2.$ $m + q = n + p$ and $p + s = q + r$

$\langle 2 \rangle 3.$ $m + q + s = n + q + r$

$\langle 2 \rangle 4.$ $m + s = n + r$

PROOF: By cancellation.

□

Definition 15.1.2 (Addition). Define *addition* $+$ on \mathbb{Z} by $[(m, n)] + [(p, q)] = [(m + p, n + q)]$.

We prove this is well-defined.

PROOF: If $m + n' = n + m'$ and $p + q' = q + p'$ then $m + p + n' + q' = n + q + m' + p'$.

□

Proposition 15.1.3. *Addition on \mathbb{Z} is commutative.*

PROOF: $[(m, n)] + [(p, q)] = [(m + p, n + q)] = [(p + m, q + n)] = [(p, q)] + [(m, n)]$.

□

Proposition 15.1.4. *Addition on \mathbb{Z} is associative.*

PROOF: $[(m, n)] + [(p, q)] + [(r, s)] = [(m + p + r, n + q + s)] = [(m, n)] + [(p, q)] + [(r, s)]$. \square

Proposition 15.1.5. *Given natural numbers m and n , we have $[(m, 0)] = [(n, 0)]$ iff $m = n$.*

PROOF: Immediate from definitions. \square

Definition 15.1.6. We identify any natural number n with the integer $[(n, 0)]$.

Proposition 15.1.7. *Addition on integers agrees with addition on natural numbers.*

PROOF: Since $[(m, 0)] + [(n, 0)] = [(m + n, 0)]$. \square

Proposition 15.1.8. *For all $a \in \mathbb{Z}$ we have $a + 0 = a$.*

PROOF: $[(m, n)] + [(0, 0)] = [(m + 0, n + 0)] = [(m, n)]$. \square

Proposition 15.1.9. *For all $a \in \mathbb{Z}$, there exists $b \in \mathbb{Z}$ such that $a + b = 0$.*

PROOF: $[(m, n)] + [(n, m)] = [(m + n, m + n)] = [(0, 0)]$ \square

Proposition 15.1.10. *The integers form an Abelian group under addition.*

PROOF: Proposition 15.1.3, 15.1.4, 15.1.8, 15.1.9. \square

Definition 15.1.11. Define multiplication \cdot on \mathbb{Z} by: $[(m, n)][(p, q)] = [(mp + nq, mq + np)]$.

We prove this is well defined.

PROOF:

$\langle 1 \rangle 1$. ASSUME: $m + n' = n + m'$ and $p + q' = q + p'$

PROVE: $mp + nq + m'q' + n'p' = mq + np + m'p' + n'q'$

$\langle 1 \rangle 2$. $mp + n'p = np + m'p$

$\langle 1 \rangle 3$. $nq + m'q = mq + n'q$

$\langle 1 \rangle 4$. $m'p + m'q' = m'q + m'p'$

$\langle 1 \rangle 5$. $n'q + n'p' = n'p + n'q'$

$\langle 1 \rangle 6$. $mp + n'p + nq + m'q + m'p + m'q' + n'q + n'p' = np + m'p + mq + n'q + m'q + m'p' + n'p + n'q'$

$\langle 1 \rangle 7$. $mp + nq + m'q' + n'p' = mq + np + m'p' + n'q'$

PROOF: By cancellation.

\square

Proposition 15.1.12. *Multiplication on integers agrees with multiplication on natural numbers.*

PROOF: Since $[(m, 0)][(n, 0)] = [(mn + 0, m0 + n0)] = [(mn, 0)]$. \square

Proposition 15.1.13. *Multiplication on \mathbb{Z} is commutative.*

PROOF: $[(m, n)][(p, q)] = [(mp + nq, mq + np)] = [(pm + qn, pn + qm)] = [(p, q)][(m, n)]$. \square

Proposition 15.1.14. *Multiplication on \mathbb{Z} is associative.*

PROOF:

$$\begin{aligned}
 [(m, n)][(p, q)][(r, s)] &= [(m, n)][(pr + qs, ps + qr)] \\
 &= [(mpr + mqs + nps + nqr, mps + mqr + npr + nqs)] \\
 &= [(mp + nq, mq + np)][(r, s)] \\
 &= [(m, n)][(p, q)][(r, s)] \quad \square
 \end{aligned}$$

Proposition 15.1.15. *Multiplication distributes over addition.*

PROOF:

$$\begin{aligned}
 [(m, n)][(p, q)] + [(m, n)][(r, s)] &= [(m, n)][(p + r, q + s)] \\
 &= [(mp + mr + nq + ns, mp + nr + mq + ms)] \\
 [(m, n)][(p, q)] + [(m, n)][(r, s)] &= [(mp + nq, mq + np)] + [(mr + ns, ms + nr)] \\
 &= [(mp + nq + mr + ns, mq + np + ms + nr)] \quad \square
 \end{aligned}$$

Proposition 15.1.16. *For any integer a we have $a1 = a$.*

PROOF: Since $[(m, n)][(1, 0)] = [(m1 + n0, m0 + n1)] = [(m, n)]$. \square

Proposition 15.1.17. *For any integers a and b , if $ab = 0$ then $a = 0$ or $b = 0$.*

PROOF:

$\langle 1 \rangle 1$. ASSUME: $[(m, n)][(p, q)] = [(0, 0)]$

$\langle 1 \rangle 2$. $mp + nq = mq + np$

$\langle 1 \rangle 3$. ASSUME: $[(m, n)] \neq [(0, 0)]$

$\langle 1 \rangle 4$. $m \neq n$

PROVE: $p = q$

$\langle 1 \rangle 5$. CASE: $m < n$

$\langle 2 \rangle 1$. $p \not\leq q$

PROOF: If $p < q$ then $mq + np < mp + nq$ by Proposition 8.4.6.

$\langle 2 \rangle 2$. $q \not\leq p$

PROOF: If $q < p$ then $mp + nq < mq + np$ by Proposition 8.4.6.

$\langle 2 \rangle 3$. $p = q$

PROOF: By trichotomy.

$\langle 1 \rangle 6$. CASE: $n < m$

PROOF: Similar.

\square

Proposition 15.1.18. *The integers \mathbb{Z} form an integral domain.*

PROOF: Propositions 15.1.13, 15.1.14, 15.1.15, 15.1.16, 15.1.17, 15.1.10. \square

Definition 15.1.19. Define $<$ on \mathbb{Z} by $[(m, n)] < [(p, q)]$ if and only if $m + q < n + p$.

We prove this is well-defined.

PROOF:

$\langle 1 \rangle 1$. ASSUME: $m + n' = n + m'$ and $p + q' = q + p'$.

PROVE: $m + q < n + p$ if and only if $m' + q' < n' + p'$

$\langle 1 \rangle 2$. $m + q < n + p$ if and only if $m' + q' < n' + p'$

PROOF:

$$m + q < n + p \Leftrightarrow m + n' + q < n + n' + p \quad (\text{Corollary 6.5.7.1})$$

$$\Leftrightarrow m' + n + q < n + n' + p$$

$$\Leftrightarrow m' + q < n' + p \quad (\text{Corollary 6.5.7.1})$$

$$\Leftrightarrow m' + q + p' < n' + p + p' \quad (\text{Corollary 6.5.7.1})$$

$$\Leftrightarrow m' + q' + p < n' + p + p'$$

$$\Leftrightarrow m' + q' < n' + p' \quad (\text{Corollary 6.5.7.1}) \square$$

Proposition 15.1.20. *The ordering on the integers agrees with the ordering on the natural numbers.*

PROOF: We have $[(m, 0)] < [(n, 0)]$ iff $m < n$. \square

Proposition 15.1.21. *$<$ is a linear order on \mathbb{Z} .*

PROOF:

$\langle 1 \rangle 1$. $<$ is irreflexive.

PROOF: We never have $m + n < m + n$.

$\langle 1 \rangle 2$. $<$ is transitive.

$\langle 2 \rangle 1$. ASSUME: $[(m, n)] < [(p, q)] < [(r, s)]$

$\langle 2 \rangle 2$. $m + q < n + p$ and $p + s < q + r$

$\langle 2 \rangle 3$. $m + q + s < n + q + r$

PROOF: $m + q + s < n + p + s < n + q + r$

$\langle 2 \rangle 4$. $m + s < n + r$

PROOF: Corollary 6.5.7.1.

$\langle 1 \rangle 3$. $<$ is total.

PROOF: Given natural numbers m, n, p and q , either $m + q < n + p$, or $m + q = n + p$, or $n + p < m + q$.

\square

Definition 15.1.22 (Positive). An integer a is *positive* iff $a > 0$.

Theorem 15.1.23. *For any integers a, b and c , we have $a < b$ if and only if $a + c < b + c$.*

PROOF:

$\langle 1 \rangle 1$. If $a < b$ then $a + c < b + c$.

$\langle 2 \rangle 1$. LET: $a = [(m, n)]$, $b = [(p, q)]$ and $c = [(r, s)]$.

$\langle 2 \rangle 2$. ASSUME: $a < b$

$\langle 2 \rangle 3$. $m + q < n + p$

$\langle 2 \rangle 4$. $m + r + q + s < n + r + p + s$

$\langle 2 \rangle 5$. $[(m + r, n + s)] < [(p + r, q + s)]$

$\langle 2 \rangle 6$. $a + c < b + c$

$\langle 1 \rangle 2$. If $a + c < b + c$ then $a < b$.

PROOF: From $\langle 1 \rangle 1$ and Proposition 5.2.6.

□

Proposition 15.1.24. *Let a , b and c be integers. If $0 < c$, then $a < b$ if and only if $ac < bc$.*

PROOF:

$\langle 1 \rangle 1$. LET: $c = [(r, s)]$

$\langle 1 \rangle 2$. ASSUME: $0 < c$

$\langle 1 \rangle 3$. $s < r$

$\langle 1 \rangle 4$. For all integers a and b , if $a < b$ then $ac < bc$

$\langle 2 \rangle 1$. LET: $a = [(m, n)]$, $b = [(p, q)]$.

$\langle 2 \rangle 2$. ASSUME: $a < b$

$\langle 2 \rangle 3$. $m + q < n + p$

$\langle 2 \rangle 4$. $(m + q)r + (p + n)s < (m + q)s + (p + n)r$

PROOF: Proposition 8.4.6, $\langle 1 \rangle 3$, $\langle 2 \rangle 3$.

$\langle 2 \rangle 5$. $mr + ns + ps + qr < ms + nr + pr + qs$

$\langle 2 \rangle 6$. $[(mr + ns, ms + nr)] < [(pr + qs, ps + qr)]$

$\langle 2 \rangle 7$. $ac < bc$

$\langle 1 \rangle 5$. For all integers a and b , if $ac < bc$ then $a < b$

PROOF: From $\langle 1 \rangle 4$ and Proposition 5.2.6.

□

Proposition 15.1.25. *Let a be a positive integer. For any integer b , there exists $k \in \mathbb{N}$ such that $b < ak$.*

PROOF:

$\langle 1 \rangle 1$. CASE: $b \leq 0$

PROOF: Take $k = 1$.

$\langle 1 \rangle 2$. CASE: $b > 0$

PROOF: Take $k = b + 1$.

□

15.2 The Rationals

Definition 15.2.1 (Rational Numbers). The set \mathbb{Q} of *rational numbers* is the field of fractions over the integers.

Proposition 15.2.2. *For any integers a and b , we have $[(a, 1)] = [(b, 1)]$ iff $a = b$.*

PROOF: Immediate from definitions. □

Henceforth we identify any integer a with the rational number $[(a, 1)]$.

Proposition 15.2.3. *Addition on the rationals agrees with addition on the integers.*

PROOF: $[(a, 1)] + [(b, 1)] = [(a \cdot 1 + b \cdot 1, 1 \cdot 1)] = [(a + b, 1)]$. \square

Proposition 15.2.4. *Multiplication on the rationals agrees with multiplication on the integers.*

PROOF: $[(a, 1)][(b, 1)] = [(ab, 1)]$ \square

Definition 15.2.5. Define the ordering $<$ on the rationals by: if b and d are positive, then $[(a, b)] < [(c, d)]$ iff $ad < bc$.

We prove this is well-defined.

PROOF:

$\langle 1 \rangle 1$. For any rational q , there exist integers a, b with b positive such that $q = [(a, b)]$.

PROOF: Since $[(a, b)] = [(-a, -b)]$, and if $b \neq 0$ then one of b and $-b$ is positive.

$\langle 1 \rangle 2$. If b, b', d and d' are positive, $[(a, b)] = [(a', b')]$, and $[(c, d)] = [(c', d')]$, then $ad < bc$ iff $a'd' < b'c'$.

PROOF:

$\langle 2 \rangle 1$. If $ad < bc$ then $a'd' < b'c'$.

$\langle 3 \rangle 1$. ASSUME: $ad < bc$

$\langle 3 \rangle 2$. $ab'd < bb'c$

$\langle 3 \rangle 3$. $a'bd < bb'c$

$\langle 3 \rangle 4$. $a'd < b'c$

$\langle 3 \rangle 5$. $a'dd' < b'cd'$

$\langle 3 \rangle 6$. $a'dd' < b'c'd$

$\langle 3 \rangle 7$. $a'd' < b'c'$

$\langle 2 \rangle 2$. If $a'd' < b'c'$ then $ad < bc$.

PROOF: Similar.

\square

Proposition 15.2.6. *The ordering on the rationals agrees with the ordering on the integers.*

PROOF: We have $[(a, 1)] < [(b, 1)]$ if and only if $a < b$. \square

Proposition 15.2.7. *The relation $<$ is a linear ordering on \mathbb{Q} .*

PROOF:

$\langle 1 \rangle 1$. $<$ is irreflexive.

PROOF: We never have $ab < ab$.

$\langle 1 \rangle 2$. $<$ is transitive.

$\langle 2 \rangle 1$. ASSUME: $[(a, b)] < [(c, d)] < [(e, f)]$ where b, d and f are positive.

$\langle 2 \rangle 2$. $ad < bc$ and $cf < de$

$\langle 2 \rangle 3$. $adf < bde$

PROOF: $adf < bcf < bde$

$\langle 2 \rangle 4$. $af < be$

$\langle 1 \rangle 3$. $<$ is total.

PROOF: For any integers a, b, c, d , we have $ad < bc$ or $ad = bc$ or $bc < ad$.

□

Proposition 15.2.8. *For any rationals r , s and t , we have $r < s$ if and only if $r + t < s + t$.*

PROOF:

⟨1⟩1. LET: a, b, c, d, e, f be integers with b, d and f positive.

⟨1⟩2. $[(a, b)] + [(e, f)] < [(c, d)] + [(e, f)]$ if and only if $[(a, b)] < [(c, d)]$.

PROOF:

$$\begin{aligned} [(a, b)] + [(e, f)] < [(c, d)] + [(e, f)] &\Leftrightarrow [(af + be, bf)] < [(cf + de, df)] \\ &\Leftrightarrow (af + be)df < (cf + de)bf \\ &\Leftrightarrow afd f + bedf < cfbf + debf \\ &\Leftrightarrow afd f < cfbf \\ &\Leftrightarrow ad < bc \\ &\Leftrightarrow [(a, b)] < [(c, d)] \end{aligned}$$

□

Corollary 15.2.8.1. *For any rational r , we have $r < 0$ if and only if $0 < -r$.*

Definition 15.2.9 (Absolute Value). For any rational r , the *absolute value* of r is defined by

$$|r| := \begin{cases} -r & \text{if } 0 < -r \\ r & \text{otherwise} \end{cases}$$

Proposition 15.2.10. *For any rationals r , s and t , if t is positive then $r < s$ iff $rt < st$.*

PROOF:

⟨1⟩1. LET: $r = [(a, b)]$, $s = [(c, d)]$ and $t = [(e, f)]$ where b, d and f are positive.

⟨1⟩2. ASSUME: $0 < t$

⟨1⟩3. $e > 0$

⟨1⟩4. $rt < st$ iff $r < s$

PROOF:

$$\begin{aligned} rt < st &\Leftrightarrow [(ae, bf)] < [(ce, df)] \\ &\Leftrightarrow aedf < cebf \\ &\Leftrightarrow ad < bc \\ &\Leftrightarrow r < s \end{aligned}$$

□

Corollary 15.2.10.1. *The rationals form an ordered field.*

Proposition 15.2.11. *Let p be a positive rational. For any rational number r , there exists $k \in \mathbb{N}$ such that $r < pk$.*

PROOF:

⟨1⟩1. LET: $p = a/b$ and $r = c/d$ where a, b and d are positive.

⟨1⟩2. PICK $k \in \mathbb{N}$ such that $bc < adk$

PROOF: Proposition 15.1.25.

⟨1⟩3. $r < pk$

□

Proposition 15.2.12. $\mathbb{Q} \approx \mathbb{N}$

PROOF: Arrange the rationals in order $0/1, 1/1, 1/2, 0/2, -1/2, -1/1, -2/1, -2/2, -2/3, -1/3, 0/3, 1/3, 2/3$, etc. then remove all duplicates. □

15.3 The Real Numbers

Definition 15.3.1 (Cauchy Sequence). A *Cauchy sequence* is a sequence (q_n) of rationals such that, for every positive rational ϵ , there exists $k \in \mathbb{N}$ such that $\forall m, n > k. |q_m - q_n| < \epsilon$.

Definition 15.3.2 (Dedekind Cut). A *Dedekind cut* is a set $x \subseteq \mathbb{Q}$ such that:

1. $\emptyset \neq x \neq \mathbb{Q}$
2. x is closed downwards.
3. x has no greatest member.

The set \mathbb{R} of *real numbers* is the set of Dedekind cuts.

Proposition 15.3.3. For any rational q , we have $\{r \in \mathbb{Q} \mid r < q\} \in \mathbb{R}$.

PROOF:

⟨1⟩1. LET: $q \in \mathbb{Q}$

⟨1⟩2. LET: $q \downarrow = \{r \mid r < q\}$

⟨1⟩3. $q \notin q \downarrow$

PROOF: We have $q - 1 \in q \downarrow$.

⟨1⟩4. $q \notin q \downarrow$

PROOF: Since $q \notin q \downarrow$.

⟨1⟩5. $q \downarrow$ is closed downwards.

PROOF: Trivial.

⟨1⟩6. $q \downarrow$ has no greatest element.

PROOF: For all $r \in q \downarrow$ we have $r < (q + r)/2 \in q \downarrow$.

□

Proposition 15.3.4. For rationals q and r , we have $q = r$ if and only if $\{s \in \mathbb{Q} \mid s < q\} = \{s \in \mathbb{Q} \mid s < r\}$.

PROOF:

⟨1⟩1. LET: $q \downarrow = \{s \in \mathbb{Q} \mid s < q\}$

⟨1⟩2. LET: $r \downarrow = \{s \in \mathbb{Q} \mid s < r\}$

⟨1⟩3. If $q = r$ then $q \downarrow = r \downarrow$

PROOF: Trivial.

$\langle 1 \rangle 4$. If $q < r$ then $q \downarrow \neq r \downarrow$

PROOF: We have $q \in r \downarrow$ and $q \notin q \downarrow$.

$\langle 1 \rangle 5$. If $r < q$ then $q \downarrow \neq r \downarrow$

PROOF: We have $r \in q \downarrow$ and $q \notin q \downarrow$.

□

Henceforth we identify a rational q with the real number $\{r \in \mathbb{Q} \mid r < q\}$.

Definition 15.3.5. Define the ordering $<$ on \mathbb{R} by: $x < y$ iff $x \subsetneq y$.

Proposition 15.3.6. *The ordering on the reals agrees with the ordering on the rationals.*

PROOF:

$\langle 1 \rangle 1$. LET: $q, r \in \mathbb{Q}$

$\langle 1 \rangle 2$. LET: $q \downarrow = \{s \in \mathbb{Q} \mid s < q\}$.

$\langle 1 \rangle 3$. LET: $r \downarrow = \{s \in \mathbb{Q} \mid s < r\}$.

PROVE: $q < r$ iff $q \downarrow \subsetneq r \downarrow$

$\langle 1 \rangle 4$. If $q < r$ then $q \downarrow \subsetneq r \downarrow$

$\langle 2 \rangle 1$. ASSUME: $q < r$

$\langle 2 \rangle 2$. $q \downarrow \subseteq r \downarrow$

PROOF: If $s < q$ then $s < r$.

$\langle 2 \rangle 3$. $q \downarrow \neq r \downarrow$

PROOF: Proposition 15.3.4.

$\langle 1 \rangle 5$. If $q \downarrow \subsetneq r \downarrow$ then $q < r$

$\langle 2 \rangle 1$. ASSUME: $q \downarrow \subsetneq r \downarrow$

$\langle 2 \rangle 2$. PICK $s \in r \downarrow$ such that $s \notin q \downarrow$

$\langle 2 \rangle 3$. $q \leq s < r$

□

Proposition 15.3.7. *The ordering $<$ is a linear ordering on \mathbb{R} .*

PROOF:

$\langle 1 \rangle 1$. $<$ is irreflexive.

PROOF: No set is a proper subset of itself.

$\langle 1 \rangle 2$. $<$ is transitive.

PROOF: Since the relationship \subsetneq is transitive on the class of all sets.

$\langle 1 \rangle 3$. $<$ is total.

$\langle 2 \rangle 1$. LET: x, y be Dedekind cuts.

$\langle 2 \rangle 2$. ASSUME: $x \not\subseteq y$

PROVE: $y \subsetneq x$

$\langle 2 \rangle 3$. PICK $q \in x$ such that $q \notin y$

$\langle 2 \rangle 4$. LET: $r \in y$

PROVE: $r \in x$

$\langle 2 \rangle 5$. $q \not\leq r$

PROOF: Since y is closed downwards.

$\langle 2 \rangle 6$. $r < q$

$\langle 2 \rangle 7$. $r \in x$

PROOF: Since x is closed downwards.

□

Proposition 15.3.8. *Any bounded nonempty subset of \mathbb{R} has a least upper bound.*

PROOF:

⟨1⟩1. LET: A be a bounded nonempty subset of \mathbb{R} .

⟨1⟩2. $\bigcup A$ is a Dedekind cut.

⟨2⟩1. $\bigcup A \neq \emptyset$

⟨3⟩1. PICK $x \in A$

⟨3⟩2. PICK $q \in x$

⟨3⟩3. $q \in \bigcup A$

⟨2⟩2. $\bigcup A \neq \mathbb{Q}$

⟨3⟩1. PICK an upper bound u for A

⟨3⟩2. PICK $q \notin u$

PROVE: $q \notin \bigcup A$

⟨3⟩3. ASSUME: for a contradiction $q \in \bigcup A$

⟨3⟩4. PICK $x \in A$ such that $q \in x$

⟨3⟩5. $x \leq u$

⟨3⟩6. $q \in u$

⟨3⟩7. Q.E.D.

PROOF: This is a contradiction.

⟨2⟩3. $\bigcup A$ is closed downwards.

⟨3⟩1. LET: $q \in \bigcup A$ and $r < q$

⟨3⟩2. PICK $x \in A$ such that $q \in x$

⟨3⟩3. $r \in x$

⟨3⟩4. $r \in \bigcup A$

⟨2⟩4. $\bigcup A$ has no greatest element.

⟨3⟩1. LET: $q \in \bigcup A$

⟨3⟩2. PICK $x \in A$ such that $q \in x$

⟨3⟩3. PICK $r \in x$ such that $q < r$

⟨3⟩4. $r \in \bigcup A$

⟨1⟩3. $\bigcup A$ is an upper bound for A .

PROOF: For all $x \in A$ we have $x \subseteq \bigcup A$.

⟨1⟩4. For any upper bound u for $\bigcup A$ we have $\bigcup A \leq u$.

PROOF: If $\forall x \in A. x \subseteq u$ we have $\bigcup A \subseteq u$.

□

Definition 15.3.9 (Addition). Define *addition* $+$ on the reals by

$$x + y := \{q + r \mid q \in x, r \in y\} .$$

We prove this is well-defined.

PROOF:

⟨1⟩1. LET: $x, y \in \mathbb{R}$

PROVE: $X + y$ is a Dedekind cut.

$\langle 1 \rangle 2. x + y \neq \emptyset$

PROOF: Pick $q \in x$ and $r \in y$; then $q + r \in x + y$.

$\langle 1 \rangle 3. x + y \neq \mathbb{Q}$

$\langle 2 \rangle 1. \text{ PICK } q \notin x \text{ and } r \notin y$

PROVE: $q + r \notin x + y$

$\langle 2 \rangle 2. \text{ ASSUME: for a contradiction } q + r \in x + y$

$\langle 2 \rangle 3. \text{ PICK } q' \in x \text{ and } r' \in y \text{ such that } q + r = q' + r'$

$\langle 2 \rangle 4. q' < q \text{ and } r' < r$

$\langle 2 \rangle 5. q' + r' < q + r$

$\langle 2 \rangle 6. \text{ Q.E.D.}$

PROOF: This is a contradiction.

$\langle 1 \rangle 4. x + y$ is closed downwards.

$\langle 2 \rangle 1. \text{ LET: } q \in x \text{ and } r \in y$

$\langle 2 \rangle 2. \text{ LET: } s < q + r$

PROVE: $s \in x + y$

$\langle 2 \rangle 3. s - r < q$

$\langle 2 \rangle 4. s - r \in x$

$\langle 2 \rangle 5. s = (s - r) + r \in x + y$

$\langle 1 \rangle 5. x + y$ has no greatest element.

$\langle 2 \rangle 1. \text{ LET: } q \in x \text{ and } r \in y$

PROVE: There exists $s \in x + y$ such that $q + r < s$

$\langle 2 \rangle 2. \text{ PICK } q' \in x \text{ and } r' \in y \text{ such that } q < q' \text{ and } r < r'$

$\langle 2 \rangle 3. q + r < q' + r' \in x + y$

□

Proposition 15.3.10. *Addition on the reals agrees with addition on the rationals.*

PROOF:

$\langle 1 \rangle 1. \text{ LET: } q, r \in \mathbb{Q}$

$\langle 1 \rangle 2. q \downarrow + r \downarrow \subseteq (q + r) \downarrow$

PROOF: If $s_1 < q$ and $s_2 < r$ then $s_1 + s_2 < q + r$.

$\langle 1 \rangle 3. (q + r) \downarrow \subseteq q \downarrow + r \downarrow$

$\langle 2 \rangle 1. \text{ LET: } s < q + r$

$\langle 2 \rangle 2. s - r < q$

$\langle 2 \rangle 3. \text{ PICK } t \text{ such that } s - r < t < q$

$\langle 2 \rangle 4. s - t < r$

$\langle 2 \rangle 5. s = t + (s - t) \in q \downarrow + r \downarrow$

□

Proposition 15.3.11. *Addition is associative.*

PROOF:

$$\begin{aligned} x + (y + z) &= \{q + r \mid q \in x, r \in y + z\} \\ &= \{q + s_1 + s_2 \mid q \in x, s_1 \in y, s_2 \in z\} \\ &= \{r + s_2 \mid r \in x + y, s_2 \in z\} \\ &= (x + y) + z \end{aligned}$$

□

Proposition 15.3.12. *Addition is commutative.*

PROOF:

$$\begin{aligned} x + y &= \{q + r \mid q \in x, r \in y\} \\ &= \{r + q \mid r \in y, q \in x\} \\ &= y + x \end{aligned}$$

□

Proposition 15.3.13. *For any $x \in \mathbb{R}$ we have $x + 0 = x$.*

PROOF:

⟨1⟩1. $x + 0 \subseteq x$

PROOF: If $q \in x$ and $r < 0$ then $q + r < q$ so $q + r \in x$.

⟨1⟩2. $x \subseteq x + 0$

⟨2⟩1. LET: $q \in x$

⟨2⟩2. PICK $r \in x$ such that $q < r$.

PROOF: x has no greatest element.

⟨2⟩3. $q - r < 0$

⟨2⟩4. $q = r + (q - r) \in x + 0$

□

Definition 15.3.14. For $x \in \mathbb{R}$, define $-x := \{q \in \mathbb{Q} \mid \exists r > q. -r \notin x\}$.

Proposition 15.3.15. *For all $x \in \mathbb{R}$ we have $-x \in \mathbb{R}$.*

PROOF:

⟨1⟩1. LET: $x \in \mathbb{R}$

⟨1⟩2. $-x \neq \emptyset$

⟨2⟩1. PICK $s \notin x$

⟨2⟩2. $-s - 1 \in -x$

⟨1⟩3. $-x \neq \mathbb{Q}$

⟨2⟩1. PICK $s \in x$

PROVE: $-s \notin -x$

⟨2⟩2. ASSUME: for a contradiction $-s \in -x$

⟨2⟩3. PICK $r > -s$ such that $-r \notin x$

⟨2⟩4. $-r < s$

⟨2⟩5. Q.E.D.

PROOF: This contradicts the fact that x is closed downwards.

⟨1⟩4. $-x$ is closed downwards.

PROOF: Immediate from definition.

⟨1⟩5. $-x$ has no greatest element.

⟨2⟩1. LET: $q \in -x$

⟨2⟩2. PICK $r > q$ such that $-r \notin x$

⟨2⟩3. PICK s such that $q < s < r$

⟨2⟩4. $s \in -x$

□

Lemma 15.3.16. *Let p be a positive rational number. For any real number x , there exists a rational $q \in x$ such that $p + q \notin x$.*

PROOF:

- $\langle 1 \rangle 1$. PICK $q_0 \in x$
- $\langle 1 \rangle 2$. There exists $k \in \mathbb{N}$ such that $q_0 + kp \notin x$
 - $\langle 2 \rangle 1$. PICK $q_1 \notin x$
 - $\langle 2 \rangle 2$. PICK $k \in \mathbb{N}$ such that $q_1 - q_0 < pk$
 - PROOF: Proposition 15.2.11.
 - $\langle 2 \rangle 3$. $q_1 < q_0 + kp$
 - $\langle 2 \rangle 4$. $q_0 + kp \notin x$
- $\langle 1 \rangle 3$. LET: k be the least natural number such that $q_0 + kp \notin x$
- $\langle 1 \rangle 4$. $k \neq 0$
- PROOF: $\langle 1 \rangle 1$
- $\langle 1 \rangle 5$. LET: $q = q_0 + (k-1)p$
- $\langle 1 \rangle 6$. $q \in x$ and $q + p \notin x$.

□

Proposition 15.3.17. *For every real x we have $x + (-x) = 0$.*

PROOF:

- $\langle 1 \rangle 1$. LET: x be a real number.
- $\langle 1 \rangle 2$. $x + (-x) \subseteq 0$
 - $\langle 2 \rangle 1$. LET: $q_1 \in x$ and $q_2 \in -x$
 - $\langle 2 \rangle 2$. PICK $r > q_2$ such that $-r \notin x$
 - $\langle 2 \rangle 3$. $q_1 < -r$
 - $\langle 2 \rangle 4$. $r < -q_1$
 - $\langle 2 \rangle 5$. $q_2 < -q_1$
 - $\langle 2 \rangle 6$. $q_1 + q_2 < 0$
- $\langle 1 \rangle 3$. $0 \subseteq x + (-x)$
 - $\langle 2 \rangle 1$. LET: $p < 0$
 - $\langle 2 \rangle 2$. $0 < -p$
 - $\langle 2 \rangle 3$. PICK $q \in x$ such that $q - p/2 \notin x$
 - PROOF: Lemma 15.3.16.
 - $\langle 2 \rangle 4$. LET: $s = p/2 - q$
 - $\langle 2 \rangle 5$. $-s \notin x$
 - $\langle 2 \rangle 6$. $p - q < s$
 - $\langle 2 \rangle 7$. $p - q \in -x$
 - $\langle 2 \rangle 8$. $p \in x + (-x)$

□

Corollary 15.3.17.1. *The reals form an Abelian group under addition.*

Proposition 15.3.18. *For any reals x, y and z , we have $x < y$ if and only if $x + z < y + z$.*

PROOF:

- $\langle 1 \rangle 1$. $\forall x, y, z \in \mathbb{R}. x \leq y \Rightarrow x + z \leq y + z$
 - $\langle 2 \rangle 1$. LET: $x, y, z \in \mathbb{R}$
 - $\langle 2 \rangle 2$. ASSUME: $x \leq y$
 - $\langle 2 \rangle 3$. For all $q \in x$ and $r \in z$ we have $q + r \in y + z$

⟨1⟩2. $\forall x, y, z \in \mathbb{R}. x + z = y + z \Leftrightarrow x = y$

PROOF: Proposition 12.1.4.

⟨1⟩3. $\forall x, y, z \in \mathbb{R}. x < y \Rightarrow x + z < y + z$

⟨1⟩4. Q.E.D.

PROOF: Proposition 5.2.6.

□

Definition 15.3.19 (Absolute Value). The *absolute value* of a real number x is defined to be

$$|x| = \begin{cases} x & \text{if } 0 \leq x \\ -x & \text{if } x < 0 \end{cases}$$

Definition 15.3.20 (Multiplication). Define *multiplication* \cdot on \mathbb{R} as follows:

- If x and y are non-negative then

$$xy = 0 \cup \{rs \mid 0 \leq r \in x \wedge 0 \leq s \in y\} .$$

- If x and y are both negative then

$$xy = (-x)(-y) .$$

- If one of x and y is negative and one is non-negative then

$$xy = -(|x||y|) .$$

We prove this is well-defined.

PROOF:

⟨1⟩1. LET: x and y be non-negative reals.

PROVE: xy is real.

⟨1⟩2. $xy \neq \emptyset$

PROOF: Since $-1 \in xy$.

⟨1⟩3. $xy \neq \mathbb{Q}$

⟨2⟩1. PICK $r \notin x$ and $s \notin y$

PROVE: $rs \notin xy$

⟨2⟩2. $0 \leq r$ and $0 \leq s$

PROOF: Since $0 \subseteq x$ and $0 \subseteq y$.

⟨2⟩3. ASSUME: for a contradiction $rs \in xy$

⟨2⟩4. PICK r' and s' such that $0 \leq r' \in x$, $0 \leq s' \in y$ and $rs = r's'$

⟨2⟩5. $r' < r$

⟨2⟩6. $s' < s$

⟨2⟩7. $r's' < rs$

⟨2⟩8. Q.E.D.

PROOF: This is a contradiction.

⟨1⟩4. xy is closed downwards.

⟨2⟩1. LET: $q \in xy$ and $r < q$

- ⟨2⟩2. CASE: $q \in 0$
 PROOF: Then $r < q < 0$ so $r \in xy$
- ⟨2⟩3. CASE: $q = s_1 s_2$ where $0 \leq s_1 \in x$ and $0 \leq s_2 \in y$
 - ⟨3⟩1. ASSUME: w.l.o.g. $0 \leq r$
 - ⟨3⟩2. $0 < s_1$ and $0 < s_2$
 - ⟨3⟩3. $r/s_2 < s_1$
 - ⟨3⟩4. $r/s_2 \in x$
 - ⟨3⟩5. $r = (r/s_2)s_2 \in xy$
- ⟨1⟩5. xy has no greatest element.
 - ⟨2⟩1. LET: $q \in xy$
 - ⟨2⟩2. CASE: $q \in 0$
 PROOF: $q < q/2 \in 0$
 - ⟨2⟩3. CASE: $q = rs$ where $0 \leq r \in x$ and $0 \leq s \in y$
 - ⟨3⟩1. PICK r' and s' with $r < r' \in x$ and $s < s' \in y$
 - ⟨3⟩2. $q < r's' \in xy$

□

Proposition 15.3.21. *Multiplication is commutative.*

PROOF: Immediate from definition. □

Proposition 15.3.22. *Multiplication is associative.*

PROOF:

- ⟨1⟩1. For non-negative reals x, y and z , we have $x(yz) = (xy)z$
 PROOF: It computes to $0 \cup \{qrs \mid 0 \leq q \in x, 0 \leq r \in y, 0 \leq s \in z\}$.
- ⟨1⟩2. For all reals x, y and z , we have $x(yz) = (xy)z$
 PROOF: It is equal to $|x||y||z|$ if an even number of them are negative, and $-(|x||y||z|)$ otherwise.

□

Proposition 15.3.23. *Multiplication distributes over addition.*

PROOF:

- ⟨1⟩1. For all non-negative reals x, y and z , we have $x(y + z) = xy + xz$
 - ⟨2⟩1. LET: x, y and z be non-negative reals.
 - ⟨2⟩2. $x(y + z) \subseteq xy + xz$
 - ⟨3⟩1. LET: $q \in x(y + z)$
 - ⟨3⟩2. CASE: $q < 0$
 PROOF: Then we have $q/2 \in xy$ and $q/2 \in xz$ so $q \in xy + xz$.
 - ⟨3⟩3. CASE: $q = rs$ where $0 \leq r \in x$ and $0 \leq s \in y + z$
 - ⟨4⟩1. PICK $s_1 \in y$ and $s_2 \in z$ such that $s = s_1 + s_2$
 - ⟨4⟩2. $rs_1 \in xy$
 PROOF: If $s_1 < 0$ then $rs_1 < 0$ so $rs_1 \in xy$. If $0 \leq s_1$ then we also have $rs_1 \in xy$.
 - ⟨4⟩3. $rs_2 \in xz$
 PROOF: Similar.
 - ⟨4⟩4. $q \in xy + xz$

PROOF: Since $q = rs_1 + rs_2$.

$\langle 2 \rangle 3$. $xy + xz \subseteq x(y + z)$

$\langle 3 \rangle 1$. LET: $q \in xy$ and $r \in xz$.

PROVE: $q + r \in x(y + z)$

$\langle 3 \rangle 2$. CASE: $q < 0$ and $r < 0$

PROOF: Then $q + r < 0$ so $q + r \in x(y + z)$.

$\langle 3 \rangle 3$. CASE: $q < 0$ and $r = r_1r_2$ where $0 \leq r_1 \in x$ and $0 \leq r_2 \in z$

$\langle 4 \rangle 1$. $q + r < r$

$\langle 4 \rangle 2$. $q + r \in xz$

$\langle 4 \rangle 3$. ASSUME: w.l.o.g. $0 \leq q + r$

PROOF: Otherwise $q + r \in x(y + z)$ immediately.

$\langle 4 \rangle 4$. PICK s_1, s_2 with $0 \leq s_1 \in x$, $0 \leq s_2 \in y$ and $q + r = s_1s_2$

$\langle 4 \rangle 5$. $s_2 \in y + z$

PROOF: Since $0 \in z$ so $s_2 = s_2 + 0 \in y + z$.

$\langle 4 \rangle 6$. $q + r \in x(y + z)$

$\langle 3 \rangle 4$. CASE: $q = q_1q_2$ where $0 \leq q_1 \in x$ and $0 \leq q_2 \in y$ and $r < 0$

PROOF: Similar.

$\langle 3 \rangle 5$. CASE: $q = q_1q_2$ where $0 \leq q_1 \in x$ and $0 \leq q_2 \in y$ and $r = r_1r_2$ where $0 \leq r_1 \in x$ and $0 \leq r_2 \in z$

$\langle 4 \rangle 1$. ASSUME: w.l.o.g. $q_1 \leq r_1$

$\langle 4 \rangle 2$. $q + r \leq r_1(q_2 + r_2) \in x(y + z)$

$\langle 1 \rangle 2$. For any negative real x and non-negative reals y and z , we have $x(y + z) = xy + xz$

PROOF:

$$\begin{aligned} x(y + z) &= -(-x)(y + z) = -((-x)y + (-x)z) & (\langle 1 \rangle 1) \\ &= -((-x)y) - ((-x)z) \\ &= xy + xz \end{aligned}$$

$\langle 1 \rangle 3$. For any non-negative real x and reals y and z with one negative and one non-negative, we have $x(y + z) = xy + xz$

$\langle 2 \rangle 1$. ASSUME: w.l.o.g. y is negative and z is non-negative.

$\langle 2 \rangle 2$. CASE: $0 \leq y + z$

PROOF:

$$\begin{aligned} xy + xz &= xy + x(-y + y + z) \\ &= -(x(-y)) + x(-y + y + z) \\ &= -(x(-y)) + x(-y) + x(y + z) & (\langle 1 \rangle 1) \\ &= x(y + z) \end{aligned}$$

$\langle 2 \rangle 3$. CASE: $y + z < 0$

$\langle 3 \rangle 1$. $-y - z > 0$

$\langle 3 \rangle 2$. $-y = z - y - z$

$\langle 3 \rangle 3$. $xy + xz = x(y + z)$

PROOF:

$$\begin{aligned}
 xy + xz &= -(x(-y)) + xz \\
 &= -(x(z - y - z)) + xz \\
 &= -(xz + x(-y - z)) + xz & ((1)1) \\
 &= -xy - x(-y - z) + xz \\
 &= -x(-y - z) \\
 &= x(y + z)
 \end{aligned}$$

(1)4. For any non-negative real x and negative reals y and z , we have $x(y + z) = xy + xz$

PROOF:

$$\begin{aligned}
 x(y + z) &= -x(-y - z) \\
 &= -(x(-y) + x(-z)) & ((1)1) \\
 &= -x(-y) - x(-z) \\
 &= xy + xz
 \end{aligned}$$

(1)5. For any negative real x and reals y and z with one negative and one non-negative, we have $x(y + z) = xy + xz$

(2)1. ASSUME: w.l.o.g. y is negative and z is non-negative.

(2)2. CASE: $0 \leq y + z$

PROOF:

$$\begin{aligned}
 x(y + z) &= -((-x)(y + z)) \\
 &= -((-x)y + (-x)z) & ((1)3) \\
 &= -((-x)y) - ((-x)z) \\
 &= (-x)(-y) - ((-x)z) \\
 &= xy + xz
 \end{aligned}$$

(2)3. CASE: $y + z < 0$

PROOF:

$$\begin{aligned}
 x(y + z) &= (-x)(-y - z) \\
 &= (-x)(-y) + (-x)(-z) & ((1)3) \\
 &= xy + xz
 \end{aligned}$$

(1)6. For any negative reals x , y and z , we have $x(y + z) = xy + xz$

PROOF:

$$\begin{aligned}
 x(y + z) &= (-x)(-y - z) \\
 &= (-x)(-y) + (-x)(-z) & ((1)1) \\
 &= xy + xz
 \end{aligned}$$

□

Proposition 15.3.24. For any real x we have $x1 = x$.

PROOF:

(1)1. CASE: $0 \leq x$

(2)1. $x1 \subseteq x$

(3)1. LET: $q \in x1$

- ⟨3⟩2. CASE: $q < 0$
 PROOF: Then $q \in x$ because $0 \leq x$.
 ⟨3⟩3. $q = rs$ where $0 \leq r \in x$ and $0 \leq s < 1$
 PROOF: Then $q < r$ so $q \in x$.
 ⟨2⟩2. $x \subseteq x1$
 ⟨3⟩1. LET: $q \in x$
 ⟨3⟩2. ASSUME: w.l.o.g. $0 \leq q$
 ⟨3⟩3. PICK r such that $q < r \in x$
 ⟨3⟩4. $0 \leq q/r < 1$
 ⟨3⟩5. $q = r(q/r) \in x1$
 ⟨1⟩2. CASE: $x < 0$
 PROOF:

$$\begin{aligned}
 x1 &= -((-x)1) \\
 &= -(-x) && (\langle 1 \rangle 1) \\
 &= x
 \end{aligned}$$

□

Lemma 15.3.25. *Let $x \in \mathbb{R}$ and c be a positive rational. Then there exists $a \in x$ and a non-least rational upper bound b for x such that $b - a = c$.*

PROOF:

- ⟨1⟩1. PICK $a_1 \in x$ such that if x has a rational supremum s then $a_1 > s - c$
 ⟨1⟩2. There exists a natural number n such that $a_1 + nc$ is an upper bound for x .
 ⟨2⟩1. PICK a non-least upper bound b_1 for x .
 ⟨2⟩2. PICK a natural number n such that $nc > b_1 - a_1$
 PROOF: Proposition 15.2.11.
 ⟨2⟩3. $a_1 + nc > b_1$
 ⟨2⟩4. $a_1 + nc$ is an upper bound for x .
 ⟨1⟩3. LET: k be the least natural number such that $a_1 + kc$ is an upper bound for x .
 ⟨1⟩4. $a_1 + (k-1)c \in x$
 ⟨1⟩5. $a_1 + kc$ is not the supremum of x .
 ⟨2⟩1. ASSUME: for a contradiction $a_1 + kc$ is the supremum of x .
 ⟨2⟩2. $a_1 > a_1 + (k-1)c$
 PROOF: ⟨1⟩1
 ⟨2⟩3. Q.E.D.
 PROOF: This is a contradiction.
 ⟨1⟩6. LET: $a = a_1 + (k-1)c$
 ⟨1⟩7. LET: $b = a_1 + kc$
 ⟨1⟩8. $b - a = c$

□

Proposition 15.3.26. *For any non-zero real x , there exists a real y such that $xy = 1$.*

PROOF:

- ⟨1⟩1. CASE: $0 < x$
- ⟨2⟩1. LET: $y = \{q \in \mathbb{Q} \mid q \leq 0\} \cup \{u^{-1} \mid u \text{ is an upper bound for } x \text{ but not the supremum of } x\}$
- ⟨2⟩2. y is a real number.
 - ⟨3⟩1. $y \neq \emptyset$
PROOF: Since $0 \in y$.
 - ⟨3⟩2. $y \neq \mathbb{Q}$
 - ⟨4⟩1. PICK $q \in x$ such that $0 < q$
 - ⟨4⟩2. $q^{-1} \notin y$
 - ⟨3⟩3. y is closed downwards.
 - ⟨4⟩1. LET: $q \in y$ and $r < q$
PROVE: $r \in y$
 - ⟨4⟩2. ASSUME: w.l.o.g. $0 < r$
 - ⟨4⟩3. q^{-1} is a non-least upper bound for x .
 - ⟨4⟩4. $q^{-1} < r^{-1}$
 - ⟨4⟩5. r^{-1} is a non-least upper bound for x .
 - ⟨4⟩6. $r \in y$
 - ⟨3⟩4. y has no greatest element.
 - ⟨4⟩1. LET: $q \in y$
PROVE: There exists $r \in y$ such that $q < r$
 - ⟨4⟩2. CASE: $q \leq 0$
 - ⟨5⟩1. PICK a non-least upper bound u for x .
 - ⟨5⟩2. $q < u^{-1} \in x$
 - ⟨4⟩3. CASE: $q = u^{-1}$ where u is a non-least upper bound for x .
 - ⟨5⟩1. PICK a non-least upper bound v with $v < u$
 - ⟨5⟩2. $u^{-1} < v^{-1} \in y$
- ⟨2⟩3. $0 < y$
- ⟨2⟩4. $xy \subseteq 1$
 - ⟨3⟩1. LET: $q \in xy$
 - ⟨3⟩2. ASSUME: w.l.o.g. $0 < q$
 - ⟨3⟩3. PICK $0 < r \in x$ and $0 < s \in y$ such that $q = rs$
 - ⟨3⟩4. s^{-1} is a non-least upper bound for x
 - ⟨3⟩5. $r < s^{-1}$
 - ⟨3⟩6. $rs < 1$
- ⟨2⟩5. $1 \subseteq xy$
 - ⟨3⟩1. LET: $q < 1$
PROVE: $q \in xy$
 - ⟨3⟩2. ASSUME: w.l.o.g. $0 < q$
 - ⟨3⟩3. PICK a_1 with $0 < a_1 \in x$
 - ⟨3⟩4. $(1 - q)a_1 > 0$
 - ⟨3⟩5. PICK $a \in x$ and a non-least upper bound w of x such that $w - a = (1 - q)a_1$
PROOF: Lemma 15.3.25.
 - ⟨3⟩6. $w - a < (1 - q)w$
 - ⟨3⟩7. $qw < a$
 - ⟨3⟩8. $w < a/q$
 - ⟨3⟩9. a/q is a non-least upper bound for x

1

PROOF:

$$\begin{aligned} yz &= (x + (y - x))z \\ &= xz + (y - x)z \\ &> xz \end{aligned}$$

1

Proposition 15.3.28.

PROOF: The function $f(x) = (2x - 1)/(x - x^2)$ is a bijection between $(0, 1)$ and \mathbb{R} . \square

$$|\mathbb{R}| = 2^{\aleph_0}$$

PROOF:

PROOF: The function H where $H(x)(n)$ is the n th binary digit of the binary expansion of x is an injection.

PROOF: Map f to the real number in $[0, 1/9]$ whose $n + 1$ st decimal digit is $f(n)$.

1

Proposition 15.3.30. *The set of algebraic numbers is countable.*

PROOF: There are countably many integer polynomials, each with finitely many roots. \square

Corollary 15.3.30.1. *There are uncountably many transcendental numbers.*

Proposition 15.3.31. *Let A be a set of disks in the plane, no two of which intersect. Then A is countable.*

PROOF: Every circle includes a point with rational coordinates. Define $f : \{q \in \mathbb{Q}^2 \mid \exists C \in A. q \in C\} \rightarrow A$ by $f(q) = C$ iff $q \in C$. Then f is surjective. \square

Proposition 15.3.32. *There exists an uncountable set of circles in the plane that do not intersect.*

PROOF: The set of all circles with origin O is uncountable. \square

Chapter 16

Complex Analysis

Theorem 16.0.1 (Hölder's Inequality). *Let p and q be real numbers with $p > 1$, $q > 1$ and $1/p + 1/q = 1$. If $(x_n) \in l^p$ and $(y_n) \in l^q$ then*

$$\sum_{n=0}^{\infty} |x_n y_n| \leq \left(\sum_{n=0}^{\infty} |x_n|^p \right)^{1/p} \left(\sum_{n=0}^{\infty} |y_n|^q \right)^{1/q}$$

PROOF:

- $\langle 1 \rangle 1$. LET: p and q be real numbers with $p > 1$ and $q > 1$
- $\langle 1 \rangle 2$. ASSUME: $1/p + 1/q = 1$
- $\langle 1 \rangle 3$. LET: $(x_n) \in l^p$
- $\langle 1 \rangle 4$. LET: $(y_n) \in l^q$
- $\langle 1 \rangle 5$. ASSUME: w.l.o.g. $x_0 \neq 0$ and $y_0 \neq 0$
- $\langle 1 \rangle 6$. For all $x \in [0, 1]$, we have

$$x^{1/p} \leq \frac{1}{p}x + \frac{1}{q}.$$

PROOF:

- $\langle 2 \rangle 1$. LET: $f : [0, 1] \rightarrow \mathbb{R}$ be the function

$$f(x) = \frac{1}{p}x + \frac{1}{q} - x^{1/p}.$$

- $\langle 2 \rangle 2$. $f'(x) = \frac{1}{p} - \frac{1}{p}x^{-1/q}$ for $x \in (0, 1]$
- $\langle 2 \rangle 3$. $f'(x) < 0$ for $x \in (0, 1]$
- $\langle 2 \rangle 4$. $f(1) = 1/p + 1/q - 1 = 0$
- $\langle 2 \rangle 5$. $f(x) \geq 0$ for all $x \in [0, 1]$
- $\langle 1 \rangle 7$. For all non-negative reals a and b , we have

$$ab \leq \frac{a^p}{p} + \frac{b^q}{q}.$$

- $\langle 2 \rangle 1$. LET: a and b be non-negative reals.
- $\langle 2 \rangle 2$. CASE: $a^p \leq b^q$
 - $\langle 3 \rangle 1$. $0 \leq a^p/b^q \leq 1$
 - $\langle 3 \rangle 2$.

$$ab^{-q/p} \leq \frac{1}{p} \frac{a^p}{b^q} + \frac{1}{q}$$

PROOF: Taking $x = a^p/b^q$ in $\langle 1 \rangle 6$.

$\langle 3 \rangle 3$.

$$ab^{1-q} \leq \frac{1}{p} \frac{a^p}{b^q} + \frac{1}{q}$$

PROOF: $-q/p = 1 - q$ from $\langle 1 \rangle 2$.

$\langle 3 \rangle 4$.

$$ab \leq \frac{a^p}{p} + \frac{b^q}{q}$$

$\langle 2 \rangle 3$. CASE: $b^q \leq a^p$

PROOF: Similar.

$\langle 1 \rangle 8$. For $j = 1, \dots, n$, we have

$$\frac{|x_j|}{(\sum_{k=0}^n |x_k|^p)^{1/p}} \frac{|y_j|}{(\sum_{k=0}^n |y_k|^q)^{1/q}} \leq \frac{1}{p} \frac{|x_j|^p}{\sum_{k=0}^n |x_k|^p} + \frac{1}{q} \frac{|y_j|^q}{\sum_{k=0}^n |y_k|^q}$$

PROOF: From $\langle 1 \rangle 7$ with

$$a = \frac{|x_j|}{(\sum_{k=0}^n |x_k|^p)^{1/p}} \text{ and } b = \frac{|y_j|}{(\sum_{k=0}^n |y_k|^q)^{1/q}}.$$

$\langle 1 \rangle 9$.

$$\frac{\sum_{j=0}^n |x_j| |y_j|}{(\sum_{k=0}^n |x_k|^p)^{1/p} (\sum_{k=0}^n |y_k|^q)^{1/q}} \leq 1$$

PROOF:

$$\begin{aligned} \frac{\sum_{j=0}^n |x_j| |y_j|}{(\sum_{k=0}^n |x_k|^p)^{1/p} (\sum_{k=0}^n |y_k|^q)^{1/q}} &\leq \frac{1}{p} + \frac{1}{q} \quad (\text{Taking the sum } j = 0 \text{ to } n \text{ in } \langle 1 \rangle 8) \\ &= 1 \end{aligned} \quad (\langle 1 \rangle 2)$$

$\langle 1 \rangle 10$. Q.E.D.

PROOF: Taking the limit $n \rightarrow \infty$ in $\langle 1 \rangle 9$.

□

Theorem 16.0.2 (Minkowski's Inequality). *Let p be a real number, $p \geq 1$. Let $(x_n), (y_n) \in l^p$. Then*

$$\left(\sum_{n=1}^{\infty} |x_n + y_n|^p \right)^{1/p} \leq \left(\sum_{n=1}^{\infty} |x_n|^p \right)^{1/p} + \left(\sum_{n=1}^{\infty} |y_n|^p \right)^{1/p}$$

PROOF:

$\langle 1 \rangle 1$. LET: p be a real number with $p \geq 1$

$\langle 1 \rangle 2$. ASSUME: w.l.o.g. $p > 1$

PROOF: The case $p = 1$ is just the Triangle Inequality.

$\langle 1 \rangle 3$. LET: q be the real such that $1/p + 1/q = 1$

$\langle 1 \rangle 4$.

$$\begin{aligned} \sum_{n=0}^{\infty} |x_n + y_n|^p &\leq \left(\sum_{n=0}^{\infty} |x_n|^p \right)^{1/p} \left(\sum_{n=0}^{\infty} |x_n + y_n|^{q(p-1)} \right)^{1/q} \\ &\quad + \left(\sum_{n=0}^{\infty} |y_n|^p \right)^{1/p} \left(\sum_{n=0}^{\infty} |x_n + y_n|^{q(p-1)} \right)^{1/q} \end{aligned}$$

PROOF:

$$\begin{aligned}
\sum_{n=0}^{\infty} |x_n + y_n|^p &= \sum_{n=0}^{\infty} |x_n + y_n| |x_n + y_n|^{p-1} \\
&\leq \sum_{n=0}^{\infty} |x_n| |x_n + y_n|^{p-1} + \sum_{n=0}^{\infty} |y_n| |x_n + y_n|^{p-1} \quad (\text{Triangle Inequality}) \\
&\leq \left(\sum_{n=0}^{\infty} |x_n|^p \right)^{1/p} \left(\sum_{n=0}^{\infty} |x_n + y_n|^{q(p-1)} \right)^{1/q} \\
&\quad + \left(\sum_{n=0}^{\infty} |y_n|^p \right)^{1/p} \left(\sum_{n=0}^{\infty} |x_n + y_n|^{q(p-1)} \right)^{1/q} \quad (\text{Hölder's Inequality})
\end{aligned}$$

⟨1⟩5.

$$\sum_{n=0}^{\infty} |x_n + y_n|^p \leq \left\{ \left(\sum_{n=0}^{\infty} |x_n|^p \right)^{1/p} + \left(\sum_{n=0}^{\infty} |y_n|^p \right)^{1/p} \right\} \left(\sum_{n=0}^{\infty} |x_n + y_n|^p \right)^{1/q}$$

⟨1⟩6. Q.E.D.

□

Chapter 17

Linear Algebra

17.1 Vector Spaces

Definition 17.1.1 (Vector Space). Let K be either \mathbb{R} or \mathbb{C} . A *vector space* over K consists of:

- a set V , whose elements are called *vectors*;
- an operation $+$: $V \times V \rightarrow V$, *addition*;
- an operation \cdot : $K \times V \rightarrow V$, *scalar multiplication*

such that:

- V is an Abelian group under $+$
- $\forall \alpha, \beta \in K. \forall x \in V. \alpha(\beta x) = (\alpha\beta)x$
- $\forall \alpha, \beta \in K. \forall x \in V. (\alpha + \beta)x = \alpha x + \beta x$
- $\forall \alpha \in K. \forall x, y \in V. \alpha(x + y) = \alpha x + \alpha y$
- $\forall x \in V. 1x = x$

We call the elements of K *scalars*. A *real vector space* is a vector space over \mathbb{R} , and a *complex vector space* is a vector space over \mathbb{C} .

Proposition 17.1.2. *Let K be either \mathbb{R} or \mathbb{C} . Let V be a vector space over K . For any $\lambda \in K$ we have $\lambda 0 = 0$.*

PROOF:

$$\begin{aligned}\lambda 0 &= \lambda(0 + 0) \\ &= \lambda 0 + \lambda 0 \\ \therefore 0 &= \lambda 0\end{aligned}$$

□

Proposition 17.1.3. *Let K be either \mathbb{R} or \mathbb{C} . Let V be a vector space over K . Let $\lambda \in K$ and $x \in V$. If $\lambda x = 0$ then either $\lambda = 0$ or $x = 0$.*

PROOF: If $\lambda \neq 0$ then $x = 1x = \lambda^{-1}\lambda x = \lambda^{-1}0 = 0$. \square

Proposition 17.1.4. *Let K be either \mathbb{R} or \mathbb{C} . Let V be a vector space over K . For any $x \in V$ we have $0x = 0$.*

PROOF:

$$\begin{aligned} 0x &= (0 + 0)x \\ &= 0x + 0x \\ \therefore 0 &= 0x \end{aligned} \quad \square$$

Proposition 17.1.5. *Let K be either \mathbb{R} or \mathbb{C} . Let V be a vector space over K . For any $x \in V$, we have $(-1)x = -x$.*

PROOF:

$$\begin{aligned} x + (-1)x &= 1x + (-1)x \\ &= (1 + (-1))x \\ &= 0x \\ &= 0 \\ \therefore (-1)x &= -x \end{aligned} \quad \square$$

Proposition 17.1.6. *Let K be either \mathbb{R} or \mathbb{C} . Then K is a vector space over K under addition and multiplication in K .*

PROOF: Easy. \square

Proposition 17.1.7. \mathbb{C} is a vector space over \mathbb{R} .

PROOF: Easy. \square

Proposition 17.1.8. *Let K be either \mathbb{R} or \mathbb{C} . Let $\{V_i\}_{i \in I}$ be a family of vector spaces over K . Then $\prod_{i \in I} V_i$ is a vector space under*

$$\begin{aligned} (f + g)(i) &= f(i) + g(i) & (f, g \in \prod_{i \in I} V_i, x \in X) \\ (\lambda f)(x) &= \lambda f(x) & (\lambda \in K, f \in \prod_{i \in I} V_i, x \in X) \end{aligned}$$

PROOF: Easy. \square

17.2 Subspaces

Definition 17.2.1 (Vector Subspace). Let K be either \mathbb{R} or \mathbb{C} . Let V be a vector space over K . A *vector subspace* of V is a subset $U \subseteq V$ such that, for all $\alpha, \beta \in K$ and $x, y \in U$, we have $\alpha x + \beta y \in U$.

It is a *proper* subspace iff $U \neq V$.

Proposition 17.2.2. *If U is a subspace of V then U is a vector space under the restrictions of $+$ and \cdot to U .*

PROOF: Easy. \square

Proposition 17.2.3. *V is a subspace of V .*

PROOF: Easy. \square

Proposition 17.2.4. *If U is a subspace of V and V is a subspace of W then U is a subspace of W .*

PROOF: Easy. \square

Definition 17.2.5. Let Ω be a topological space. Then $\mathcal{C}(\Omega)$ is the complex vector space of all continuous functions from Ω to \mathbb{C} . This is a subspace of \mathbb{C}^Ω .

Definition 17.2.6. Let $n, k \in \mathbb{N}$. Let Ω be an open subset of \mathbb{R}^n . Then $\mathcal{C}^k(\Omega)$ is the complex vector space of all functions $\Omega \rightarrow \mathbb{C}$ that have all continuous partial derivatives of order k . This is a subspace of $\mathcal{C}(\Omega)$. If $l > k$ then $\mathcal{C}^l(\Omega)$ is a subspace of $\mathcal{C}^k(\Omega)$.

Definition 17.2.7. Let $n \in \mathbb{N}$. Let Ω be an open subset of \mathbb{R}^n . Then $\mathcal{C}^\infty(\Omega)$ is the complex vector space of all infinitely differentiable functions $\Omega \rightarrow \mathbb{C}$. This is a subspace of $\mathcal{C}^k(\Omega)$ for all k .

Definition 17.2.8. Let $n \in \mathbb{N}$. Let Ω be an open subset of \mathbb{R}^n . Then $\mathcal{P}(\Omega)$ is the complex vector space of all complex polynomials of n variables, considered as functions $\Omega \rightarrow \mathbb{C}$. This is a subspace of $\mathcal{C}^\infty(\Omega)$.

Proposition 17.2.9. *The space of all convergent sequences in \mathbb{C} is a subspace of the space of all bounded sequences in \mathbb{C} , which is a subspace of $\mathbb{C}^\mathbb{N}$.*

PROOF: Easy. \square

Definition 17.2.10. Let p be a real number, $p \geq 1$. Let l^p be the set of all complex sequences (z_n) such that $\sum_{n=1}^\infty |z_n|^p < \infty$.

Proposition 17.2.11. *For p a real number ≥ 1 , we have that l^p is a subspace of $\mathbb{C}^\mathbb{N}$.*

PROOF:

(1)1. For all $(x_n), (y_n) \in l^p$, we have $(x_n + y_n) \in l^p$.

PROOF: From Minkowski's Inequality.

(1)2. For all $\lambda \in \mathbb{C}$ and $(x_n) \in l^p$ we have $(\lambda x_n) \in l^p$

PROOF:

$$\sum_{n=1}^\infty |\lambda x_n|^p = |\lambda|^p \sum_{n=1}^\infty |x_n|^p < \infty$$

\square

Definition 17.2.12 (Linear Combination). Let K be either \mathbb{R} or \mathbb{C} . Let V be a vector space over K . Let $x, x_1, \dots, x_n \in V$. Then x is a *linear combination* of x_1, \dots, x_n iff there exist $\alpha_1, \dots, \alpha_n \in K$ such that

$$x = \alpha_1 x_1 + \dots + \alpha_n x_n .$$

Definition 17.2.13 (Linearly Independent). A finite set of vectors $\{x_1, \dots, x_n\}$ is *linearly independent* iff, whenever $\alpha_1 x_1 + \dots + \alpha_n x_n = 0$, then $\alpha_1 = \dots = \alpha_n = 0$.

A set of vectors is *linearly independent* iff every finite subset is linearly independent; otherwise, it is *linearly dependent*.

Definition 17.2.14 (Span). Let \mathcal{A} be a set of vectors. The *span* of \mathcal{A} , $\text{span } \mathcal{A}$, is the set of all linear combinations of elements of \mathcal{A} .

Proposition 17.2.15. $\text{span } \mathcal{A}$ is the smallest subspace of V that includes \mathcal{A} .

PROOF: Easy. \square

Definition 17.2.16 (Basis). A *basis* for V is a linearly independent set of vectors \mathcal{B} such that $\text{span } \mathcal{B} = V$.

Definition 17.2.17 (Finite Dimensional). A vector space is *finite dimensional* iff it has a finite basis; otherwise it is *infinite dimensional*.

Proposition 17.2.18. In a finite dimensional vector space, any two bases have the same number of elements.

Definition 17.2.19 (Dimension). The *dimension* of a finite dimensional vector space V , $\dim V$, is the number of elements in any basis.

Proposition 17.2.20.

$$\dim K^n = n$$

PROOF: The standard basis is the set of vectors with one coordinate 1 and all others 0. \square

Proposition 17.2.21. The dimension of \mathbb{C}^n as a real vector space is $2n$.

17.3 Normed Spaces

Definition 17.3.1 (Norm). Let K be either \mathbb{R} or \mathbb{C} . A *norm* on a vector space V over K is a function $\| \cdot \| : V \rightarrow \mathbb{R}$ such that:

1. $\forall x \in V, \|x\| = 0 \Rightarrow x = 0$
2. $\forall \lambda \in K, \forall x \in V, \|\lambda x\| = |\lambda| \|x\|$
3. *Triangle Inequality* $\forall x, y \in V, \|x + y\| \leq \|x\| + \|y\|$

Proposition 17.3.2. *Let K be either \mathbb{R} or \mathbb{C} . Let V be a vector space over K . Let $\| \cdot \|$ be a norm on V . Then*

$$\|0\| = 0 \text{ .}$$

PROOF:

$$\begin{aligned} \|0\| &= \|0 \cdot 0\| && \text{(Proposition 17.1.4)} \\ &= |0|\|0\| && \text{(Axiom 2 for a norm)} \\ &= 0 && \square \end{aligned}$$

Proposition 17.3.3. *Let K be either \mathbb{R} or \mathbb{C} . Let V be a vector space over K . Let $\| \cdot \|$ be a norm on V . Let $x \in V$. Then*

$$\|x\| \geq 0 \text{ .}$$

PROOF:

$$\begin{aligned} 0 &= \|0\| && \text{(Proposition 17.3.2)} \\ &= \|x - x\| \\ &\leq \|x\| + \|-x\| && \text{(Triangle Inequality)} \\ &= \|x\| + \|x\| && \text{(Axiom 2 for a norm)} \\ &= 2\|x\| && \square \end{aligned}$$

Proposition 17.3.4. *Let K be either \mathbb{R} or \mathbb{C} . Let V be a vector space over K . Let $\| \cdot \|$ be a norm on V . Let $x, y \in V$. Then*

$$\| \|x\| - \|y\| \| \leq \|x - y\| \text{ .}$$

PROOF:

$$\langle 1 \rangle 1. \quad \|x\| - \|y\| \leq \|x - y\|$$

PROOF: $\|x\| \leq \|x - y\| + \|y\|$ by the Triangle Inequality.

$$\langle 1 \rangle 2. \quad \|y\| - \|x\| \leq \|x - y\|$$

PROOF:

$$\begin{aligned} \|x\| + \|x - y\| &= \|x\| + \|y - x\| && \text{(Axiom 2 for a norm)} \\ &\leq \|y\| && \text{(Triangle Inequality)} \end{aligned}$$

\square

Definition 17.3.5 (Euclidean Norm). The *Euclidean norm* on \mathbb{C}^n is defined by

$$\|(z_1, \dots, z_n)\| = \sqrt{|z_1|^2 + \dots + |z_n|^2}$$

Proposition 17.3.6. *Define $\| \cdot \| : \mathbb{C}^n \rightarrow \mathbb{R}$ by*

$$\|(z_1, \dots, z_n)\| = |z_1| + \dots + |z_n|$$

Then this defines a norm on \mathbb{C}^n .

PROOF: Easy. \square

Proposition 17.3.7. Define $\| \cdot \| : \mathbb{C}^n \rightarrow \mathbb{R}$ by

$$\|(z_1, \dots, z_n)\| = \max(|z_1|, \dots, |z_n|)$$

Then this defines a norm on \mathbb{C}^n .

PROOF: Easy. \square

Proposition 17.3.8. Let Ω be a closed bounded subset of \mathbb{R}^n . Define $\| \cdot \| : \mathcal{C}(\Omega) \rightarrow \mathbb{R}$ by $\|f\| = \max_{x \in \Omega} |f(x)|$. Then $\| \cdot \|$ defines a norm on $\mathcal{C}(\Omega)$.

PROOF: Easy. \square

Proposition 17.3.9. Let p be a real number, $p \geq 1$. Define $\| \cdot \| : l^p \rightarrow \mathbb{R}$ by

$$\|(z_n)\| = \left(\sum_{n=0}^{\infty} |z_n|^p \right)^{1/p}.$$

Then this defines a norm on l^p .

PROOF: Easy. The triangle inequality is Minkowski's Inequality. \square

Definition 17.3.10 (Normed Space). Let K be either \mathbb{R} or \mathbb{C} . A *normed space* over K consists of a vector space V over K and a norm on V .

We shall write simply:

- K^n for the normed space K^n under the Euclidean norm
- $\mathcal{C}(\Omega)$ for the normed space $\mathcal{C}(\Omega)$ under the norm $\|f\| = \max_{x \in \Omega} |f(x)|$, called the *uniform convergence norm*
- l^p for the normed space l^p under the norm $\|(z_n)\| = (\sum_{n=0}^{\infty} |z_n|^p)^{1/p}$.

Definition 17.3.11 (Convergence). Let K be either \mathbb{R} or \mathbb{C} . Let V be a normed space over K . Let (x_n) be a sequence in V and $l \in V$. Then x_n *converges* to l , $x_n \rightarrow l$ as $n \rightarrow \infty$, iff for all $\epsilon > 0$ there exists N such that, for all $n \geq N$, we have $\|x_n - l\| < \epsilon$.

Proposition 17.3.12. Let K be either \mathbb{R} or \mathbb{C} . Let V be a normed space over K . Let (x_n) be a sequence in V and $l \in V$. Then $x_n \rightarrow l$ as $n \rightarrow \infty$ in V if and only if $\|x_n - l\| \rightarrow 0$ as $n \rightarrow \infty$ in \mathbb{R} .

PROOF: Immediate from definitions. \square

Proposition 17.3.13. A convergent sequence in a normed space has at most one limit.

PROOF:

$\langle 1 \rangle$ 1. ASSUME: $x_n \rightarrow l$ and $x_n \rightarrow m$ as $n \rightarrow \infty$

$\langle 1 \rangle$ 2. $\|l - m\| = 0$

⟨2⟩1. LET: $\epsilon > 0$
 PROVE: $\|l - m\| < \epsilon$
 ⟨2⟩2. PICK N such that $\forall n \geq N, \|x_n - l\| < \epsilon/2$ and $\forall n \geq N, \|x_n - m\| < \epsilon/2$
 ⟨2⟩3. $\|l - m\| < \epsilon$
 PROOF: $\|l - m\| \leq \|x_N - l\| + \|x_N - m\| < \epsilon$
 ⟨1⟩3. $l = m$
 \square

Proposition 17.3.14. *Let K be either \mathbb{R} or \mathbb{C} . Let V be a normed space over K . If $x_n \rightarrow l$ as $n \rightarrow \infty$ in V and $\lambda_n \rightarrow \lambda$ as $n \rightarrow \infty$ in K , then $\lambda_n x_n \rightarrow \lambda l$ as $n \rightarrow \infty$.*

PROOF:

⟨1⟩1. LET: $\epsilon > 0$
 ⟨1⟩2. LET: $K = |\lambda| + \epsilon/2\|l\|$
 ⟨1⟩3. PICK N such that, for all $n \geq N$, we have $|\lambda_n - \lambda| < \epsilon/2\|l\|$ and $\|x_n - l\| < \epsilon/(2K)$
 ⟨1⟩4. For all $n \geq N$ we have $|\lambda_n| < K$
 ⟨1⟩5. $\|\lambda_n x_n - \lambda l\| < \epsilon$

PROOF:

$$\begin{aligned}
 \|\lambda_n x_n - \lambda l\| &\leq \|\lambda_n x_n - \lambda_n l\| + \|\lambda_n l - \lambda l\| \\
 &= |\lambda_n| \|x_n - l\| + |\lambda_n - \lambda| \|l\| \\
 &< K \frac{\epsilon}{2K} + \frac{\epsilon}{2\|l\|} \|l\| \\
 &= \epsilon
 \end{aligned}$$

\square

Proposition 17.3.15. *In a normed space, if $x_n \rightarrow l$ and $y_n \rightarrow m$ then $x_n + y_n \rightarrow l + m$*

PROOF:

$$\begin{aligned}
 \|(x_n + y_n) - (l + m)\| &\leq \|x_n - l\| + \|y_n - m\| \\
 &\rightarrow 0
 \end{aligned}$$

\square

Definition 17.3.16 (Uniform Convergence). Let Ω be a closed bounded set in \mathbb{R}^n . Let (f_n) be a sequence in $\mathcal{C}(\Omega)$ and $f \in \mathcal{C}(\Omega)$. Then (f_n) converges uniformly to f if and only if, for every $\epsilon > 0$, there exists N such that $\forall x \in \Omega, \forall n \geq N, |f_n(x) - f(x)| < \epsilon$.

Proposition 17.3.17. *(f_n) converges uniformly to f iff (f_n) converges to f under the uniform convergence norm.*

PROOF: Easy. \square