Mathematics

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August 9, 2023

Contents

Chapter 1

Sets and Classes

1.1 Classes

Our language is the language of first-order logic with equality over one primitive binary predicate $\in$. We call all the objects we reason about *sets*. When $a \in b$, we say a is a *member* or *element* of b, or b *contains* a. We write $b \ni a$ for $a \in b$, and $a \notin b$ for $\neg(a \in b)$. We write $\forall x \in a.\phi$ as an abbreviation for $\forall x(x \in a \to \phi)$, and $\exists x \in a.\phi$ as an abbreviation for $\exists x(x \in a \wedge \phi)$.

We shall speak informally of *classes* as an abbreviation for talking about predicates. A *class* is determined by a unary predicate $\phi[x]$ (possibly with parameters). We write $\{x \mid \phi[x]\}$ or $\{x : \phi[x]\}$ for the class determined by $\phi[x]$. We write 'a is an element of $\{x \mid \phi[x]\}$' or '$a \in \{x \mid \phi[x]\}$' for $\phi[a]$.

We write $\{t[x_1, \ldots, x_n] \mid P[x_1, \ldots, x_n]\}$ for

$$\{y \mid \exists x_1, \ldots, x_n(y = t[x_1, \ldots, x_n] \wedge P[x_1, \ldots, x_n])\} \ .$$

We say two classes $\mathbf{A}$ and $\mathbf{B}$ are *equal*, and write $\mathbf{A} = \mathbf{B}$, iff $\forall x(x \in \mathbf{A} \leftrightarrow x \in \mathbf{B})$.

Proposition Schema 1.1.1. *For any class $\mathbf{A}$, the following is a theorem.*

$$\mathbf{A} = \mathbf{A}$$

PROOF: We have $\forall x(x \in \mathbf{A} \Leftrightarrow x \in \mathbf{A})$. $\square$

Proposition Schema 1.1.2. *For any classes $\mathbf{A}$ and $\mathbf{B}$, the following is a theorem.*
If $\mathbf{A} = \mathbf{B}$ then $\mathbf{B} = \mathbf{A}$.

PROOF: If $\forall x(x \in \mathbf{A} \Leftrightarrow x \in \mathbf{B})$ then $\forall x(x \in \mathbf{B} \Leftrightarrow x \in \mathbf{A})$. $\square$

Proposition Schema 1.1.3. *For any classes $\mathbf{A}$, $\mathbf{B}$ and $\mathbf{C}$, the following is a theorem.*
If $\mathbf{A} = \mathbf{B}$ and $\mathbf{B} = \mathbf{C}$ then $\mathbf{A} = \mathbf{C}$.

PROOF: If $\forall x(x \in \mathbf{A} \Leftrightarrow x \in \mathbf{B})$ and $\forall x(x \in \mathbf{B} \Leftrightarrow x \in \mathbf{C})$ then $\forall x(x \in \mathbf{A} \Leftrightarrow x \in \mathbf{C})$. $\square$

1.1.1   Subclasses

Definition 1.1.4 (Subclass)**.** We say a class **A** is a *subclass* of **B**, or **B** is a *superclass* of **A**, or **B** *includes* **A**, and write **A** $\subseteq$ **B** or **B** $\supseteq$ **A**, iff every element of **A** is an element of **B**. Otherwise we write **A** $\nsubseteq$ **B** or **B** $\nsupseteq$ **A**.

We say **A** is a *proper* subclass of **B**, **B** is a *proper* superclass of **A**, or **B** *properly* includes **A**, and write **A** $\subsetneq$ **B** or **B** $\supsetneq$ **A**, iff **A** $\subseteq$ **B** and **A** $\neq$ **B**.

Proposition Schema 1.1.5. *For any class* **A***, the following is a theorem.*

$$\mathbf{A} \subseteq \mathbf{A}$$

PROOF: Every element of **A** is an element of **A**. □

Proposition Schema 1.1.6. *For any classes* **A** *and* **B***, the following is a theorem.*
If **A** $\subseteq$ **B** *and* **B** $\subseteq$ **A** *then* **A** = **B***.*

PROOF: If every element of **A** is an element of **B**, and every element of **B** is an element of **A**, then **A** and **B** have exactly the same elements. □

Proposition Schema 1.1.7. *For any classes* **A***,* **B** *and* **C***, the following is a theorem.*
If **A** $\subseteq$ **B** *and* **B** $\subseteq$ **C** *then* **A** $\subseteq$ **C***.*

PROOF: If every element of **A** is an element of **B** and every element of **B** is an element of **C** then every element of **A** is an element of **C**. □

1.1.2   Constructions of Classes

Definition 1.1.8 (Empty Class)**.** The *empty class* $\emptyset$ is $\{x \mid \bot\}$. Every other class is *nonempty.*

Definition 1.1.9 (Universal Class)**.** The *universal class* **V** is $\{x \mid \top\}$.

Definition 1.1.10 (Enumeration)**.** Given objects a_1, $\ldots$, a_n, we define the class $\{a_1, \ldots, a_n\}$ to be the class $\{x \mid x = a_1 \vee \cdots \vee x = a_n\}$.

Definition 1.1.11 (Intersection)**.** For any classes **A** and **B**, the *intersection* **A** $\cap$ **B** is $\{x \mid x \in \mathbf{A} \wedge x \in \mathbf{B}\}$.

Definition 1.1.12 (Union)**.** For any classes **A** and **B**, the *union* **A** $\cup$ **B** is $\{x \mid x \in \mathbf{A} \vee x \in \mathbf{B}\}$.

Definition 1.1.13 (Relative Complement)**.** Let **A** and **B** be classes. The *relative complement* of **B** in **A** is the class **A** $-$ **B** $:= \{x \in \mathbf{A} \mid x \notin \mathbf{B}\}$.

Definition 1.1.14 (Symmetric Difference)**.** For any classes **A** and **B**, the *symmetric difference* is the class **A** $+$ **B** $:= (\mathbf{A} - \mathbf{B}) \cup (\mathbf{B} - \mathbf{A})$.

Definition 1.1.15 (Pairwise disjoint)**.** Let **A** be a class. We say the elements of **A** are *pairwise disjoint* iff, for all $x, y \in \mathbf{A}$, if $x \cap y \neq \emptyset$ then $x = y$.

1.2 Sets and the Axiom of Extensionality

Definition 1.2.1 (Axiom of Extensionality)**.** The *Axiom of Extensionality* is the statement: if two sets have exactly the same members, then they are equal.

$$\forall x, y(\forall z(z \in x \Leftrightarrow z \in y) \Rightarrow x = y) \ .$$

When working in a theory with the Axiom of Extensionality, we may identify a set a with the class $\{x \mid x \in a\}$. Our use of the symbols $\in$ and $=$ is consistent. We say a class **A** *is a set* iff there exists a set a such that $a = \mathbf{A}$; that is, $\{x \mid \phi[x]\}$ is a set iff $\exists a \forall x(x \in a \leftrightarrow \phi[x])$. Otherwise, **A** is a *proper class*.

Definition 1.2.2 (Subset)**.** If A is a set and $A \subseteq \mathbf{B}$, we say A is a *subset* of **B**.

Definition 1.2.3 (Union)**.** The *union* of a class **A** is $\{x \mid \exists X \in \mathbf{A}.x \in X\}$. We write $\bigcup_{P(x)} t(x)$ for $\bigcup\{t(x) \mid P(x)\}$.

Definition 1.2.4 (Intersection)**.** The *intersection* of a class **A** is $\{x \mid \forall X \in \mathbf{A}.x \in X\}$. We write $\bigcap_{P(x)} t(x)$ for $\bigcap\{t(x) \mid P(x)\}$.

Definition 1.2.5 (Power Class)**.** For any class **A**, the *power class* $\mathcal{P}\mathbf{A}$ is $\{X \mid X \subseteq \mathbf{A}\}$.

1.3 The Other Axioms

Definition 1.3.1 (Pairing Axiom)**.** The *Pairing Axiom* is the statement: for any sets a and b, the class $\{a, b\}$ is a set.

$$\forall a \forall b \exists c \forall x(x \in c \Leftrightarrow x = a \vee x = b)$$

Definition 1.3.2 (Union Axiom)**.** The *Union Axiom* is the statement: for any set A, the class $\bigcup A$ is a set.

$$\forall A \exists B \forall x(x \in B \Leftrightarrow \exists y(y \in A \wedge x \in y))$$

Definition 1.3.3 (Comprehension Axiom Scheme)**.** The *Comprehension Axiom Scheme* is the set of sentences of the form, for any class **A**: If **A** is a subclass of a set then **A** is a set.

That is, for any property $P[x, y_1, \ldots, y_n]$:
For any sets a_1, ..., a_n and B, the class $\{x \in B \mid P[x, a_1, \ldots, a_n]\}$ is a set.

$$\forall a_1, \ldots, a_n, B. \exists C. \forall x(x \in C \Leftrightarrow x \in B \wedge P[x, a_1, \ldots, a_n])$$

Definition 1.3.4 (Replacement Axiom Scheme)**.** The *Replacement Axiom Scheme* is the set of sentences of the form, for some property $P[x, y, z_1, \ldots, z_n]$:

For any sets a_1, ..., a_n, B, assume for all $x \in B$ there exists at most one y such that $P[x, y, a_1, \ldots, a_n]$. Then $\{y \mid \exists x \in B.P[x, y, a_1, \ldots, a_n]\}$ is a set.

$$\forall a_1, \ldots, a_n, B(\forall x \in B.\forall y, y'(P[x, y, a_1, \ldots, a_n] \wedge P[x, y', a_1, \ldots, a_n] \Rightarrow y = y') \Rightarrow$$
$$\exists C \forall y(y \in C \Leftrightarrow \exists x \in B.P[x, y, a_1, \ldots, a_n]))$$

Definition 1.3.5 (Power Set Axiom)**.** The *Power Set Axiom* is the statement: the power class of a set is a set.

$$\forall A \exists B \forall x (x \in B \Leftrightarrow \forall y (y \in x \Rightarrow y \in A))$$

Definition 1.3.6 (Axiom of Infinity)**.** The *Axiom of Infinity* is the statement: there exists a set I such that $\emptyset \in I$ and $\forall x \in I.x \cup \{x\} \in I$.

$$\exists I (\exists e \in I.\forall x.x \notin e \wedge \forall x \in I.\exists y \in I.\forall z (z \in y \Leftrightarrow z \in x \vee z = x))$$

Definition 1.3.7 (Axiom of Choice)**.** The *Axiom of Choice* is the statement: For any set A of pairwise disjoint, nonempty sets, there exists a set C such that, for all $x \in A$, we have $x \cap C$ has exactly one element.

$$\forall A (\forall x \in A.\exists yy \in x \wedge$$
$$\forall x, y \in A.\forall z (z \in x \wedge z \in y \Rightarrow x = y) \Rightarrow$$
$$\exists C.\forall x \in A.\exists y \forall z (z \in x \wedge z \in C \Leftrightarrow z = y))$$

Definition 1.3.8 (Axiom of Regularity)**.** The *Axiom of Regularity* is the statement: for any A, if A has a member, then there exists $m \in A$ such that $m \cap A = \emptyset$.

$$\forall A (\exists x.x \in A \Rightarrow \exists m \in A.\neg \exists x (x \in m \wedge x \in A))$$

Definition 1.3.9 (Zermelo Set Theory)**.** *Zermelo set theory* is the theory whose axioms are:

- Extensionality

- Pairing

- Union

- Comprehension

- Power Set

- Infinity

- Choice

- Regularity

We label theorems with Z when they are provable in Zermelo set theory.

Definition 1.3.10 (Zermelo-Fraenkel Set Theory)**.** *Zermelo-Fraenkel set theory* is the theory whose axioms are:

- Extensionality

- Union

- Replacement

- Power Set

- Infinity

- Choice

- Regularity

We label theorems with ZFC when they are provable in Zermelo-Fraenkel set theory.

We label a theorem with FOL if it can be proved in first-order logic, i.e. from no axioms.

1.4 ZFC Extends Z

Proposition 1.4.1 (Z,ZFC). *The empty class $\emptyset$ is a set.*

PROOF: Immediate from the Axiom of Infinity. □

Proposition 1.4.2 (ZFC). *The Axiom of Pairing is a theorem of ZFC.*

PROOF:
$\langle 1 \rangle$1. LET: a, b be sets.
$\langle 1 \rangle$2. LET: $P(x, y)$ be the predicate $(x = \emptyset \wedge y = a) \vee (x = \mathcal{P}\emptyset \wedge y = b)$.
$\langle 1 \rangle$3. For all $x \in \mathcal{P}\mathcal{P}\emptyset$, there exists at most one y such that $P(x, y)$.
  $\langle 2 \rangle$1. LET: $x \in \mathcal{P}\mathcal{P}\emptyset$
  $\langle 2 \rangle$2. LET: y and y' be sets.
  $\langle 2 \rangle$3. ASSUME: $P(x, y)$ and $P(x, y')$
  $\langle 2 \rangle$4. $(x = \emptyset \wedge y = a) \vee (x = \mathcal{P}\emptyset \wedge y = b)$
    PROOF: From $\langle 2 \rangle$3.
  $\langle 2 \rangle$5. $(x = \emptyset \wedge y' = a) \vee (x = \mathcal{P}\emptyset \wedge y' = b)$
    PROOF: From $\langle 2 \rangle$3.
  $\langle 2 \rangle$6. $\emptyset \neq \mathcal{P}\emptyset$
    PROOF: Since $\emptyset \in \mathcal{P}\emptyset$ and $\emptyset \notin \emptyset$.
  $\langle 2 \rangle$7. $y = y'$
$\langle 1 \rangle$4. LET: A be the set $\{y \mid \exists x \in \mathcal{P}\mathcal{P}\emptyset . P(x, y)\}$.
$\langle 1 \rangle$5. $A = \{a, b\}$
□

Proposition Schema 1.4.3 (ZFC). *Every instance of the Comprehension Axiom Scheme is a theorem of ZFC.*

PROOF:
$\langle 1 \rangle$1. LET: $P(x)$ be a predicate.
$\langle 1 \rangle$2. LET: A be a set.
$\langle 1 \rangle$3. LET: $Q(x, y)$ be the predicate $P(x) \wedge y = x$.

⟨1⟩4. For all $x \in A$, there exists at most one y such that $Q(x, y)$.
    ⟨2⟩1. LET: $x \in A$
    ⟨2⟩2. LET: y and y' be sets.
    ⟨2⟩3. ASSUME: $Q(x, y)$ and $Q(x, y')$
    ⟨2⟩4. $x \in A \wedge P(x) \wedge y = x \wedge y' = x$
      PROOF: From ⟨2⟩3.
    ⟨2⟩5. $y = y'$
      PROOF: From ⟨2⟩4.
⟨1⟩5. LET: B be the set $\{y \mid \exists x \in A.Q(x, y)\}$
  PROOF: This is a set by an Axiom of Replacement and ⟨1⟩4.
⟨1⟩6. $B = \{y \in A \mid P(y)\}$
  PROOF:

$$
\begin{aligned}
y \in B &\Leftrightarrow \exists x \in A.Q(x, y) & (⟨1⟩5) \\
&\Leftrightarrow \exists x \in A(P(x) \wedge y = x) & (⟨1⟩3) \\
&\Leftrightarrow P(y)
\end{aligned}
$$

☐

Corollary Schema 1.4.3.1 (ZFC). *Every axiom of Z is a theorem of ZFC.*

It follows that every theorem of Z is a theorem of ZFC.

1.5   Consequences of the Axioms

Proposition 1.5.1 (Z). *The union of two sets is a set.*

PROOF: Because $A \cup B = \bigcup\{A, B\}$. ☐

Proposition Schema 1.5.2 (Z). *For any number n, the following is a theorem:*
    For any sets $a_1, \ldots, a_n$, the class $\{a_1, \ldots, a_n\} = \{x \mid x = a_1 \vee \cdots \vee x = a_n\}$
is a set.

PROOF: The case $n = 1$ follows from Pairing since $\{a\} = \{a, a\}$.
If we have proved the theorem for n we have $\{a_1, \ldots, a_n, a_{n+1}\} = \{a_1, \ldots, a_n\} \cup \{a_{n+1}\}$. ☐

Proposition 1.5.3 (Z). *No set is a member of itself.*

PROOF:
⟨1⟩1. LET: x be any set.
⟨1⟩2. PICK $m \in \{x\}$ such that $m \cap \{x\} = \emptyset$.
  PROOF: Axiom of Regularity.
⟨1⟩3. $m = x$
⟨1⟩4. $x \cap \{x\} = \emptyset$
⟨1⟩5. $x \notin x$
☐

Corollary 1.5.3.1 (Z). *The universal class $\mathbf{V}$ is a proper class.*

PROOF: If $\mathbf{V}$ is a set then $\mathbf{V} \in \mathbf{V}$, contradicting the Proposition. □

Proposition 1.5.4 (Z)**.** *There are no sets a and b such that $a \in b$ and $b \in a$.*

PROOF:
⟨1⟩1. LET: a and b be any sets.
⟨1⟩2. PICK $m \in \{a, b\}$ such that $m \cap \{a, b\} = \emptyset$
⟨1⟩3. CASE: $m = a$
  PROOF: Then $b \notin a$.
⟨1⟩4. CASE: $m = b$
  PROOF: Then $a \notin b$.
□

Proposition 1.5.5 (Z)**.** *The intersection of a set and a class is a set.*

PROOF: Immediate from Comprehension. □

Proposition 1.5.6 (Z)**.** *The relative complement of a class in a set is a set.*

  [Z]

PROOF: Immediate from Comprehension. □

Corollary 1.5.6.1 (Z)**.** *The symmetric difference of two sets is a set.*

Proposition 1.5.7 (Z)**.** *The intersection of a nonempty class is a set.*

PROOF:
⟨1⟩1. LET: $\mathbf{A}$ be a nonempty class.
⟨1⟩2. PICK $B \in \mathbf{A}$
⟨1⟩3. $\bigcap \mathbf{A} \subseteq B$
⟨1⟩4. $\bigcap \mathbf{A}$ is a set.
  PROOF: By Comprehension.
□

Proposition Schema 1.5.8 (FOL)**.** *For any classes $\mathbf{A}$ and $\mathbf{B}$, the following is a theorem:*
  If $\mathbf{A} \subseteq \mathbf{B}$ then $\mathcal{P}\mathbf{A} \subseteq \mathcal{P}\mathbf{B}$.

PROOF: Every subset of $\mathbf{A}$ is a subset of $\mathbf{B}$. □

Proposition Schema 1.5.9 (FOL)**.** *For any classes $\mathbf{A}$ and $\mathbf{B}$, the following is a theorem:*
  If $\mathbf{A} \subseteq \mathbf{B}$ then $\bigcup \mathbf{A} \subseteq \bigcup \mathbf{B}$.

PROOF: If $x \in X \in \mathbf{A}$ then $x \in X \in \mathbf{B}$. □

Proposition Schema 1.5.10 (Z)**.** *For any class $\mathbf{A}$, the following is a theorem:*

$$\mathbf{A} = \bigcup \mathcal{P}\mathbf{A}$$

PROOF:

$\langle 1 \rangle 1.$ $\mathbf{A} \subseteq \bigcup \mathcal{P} \mathbf{A}$
   PROOF: For all $x \in \mathbf{A}$ we have $x \in \{x\} \in \mathcal{P} \mathbf{A}$.
$\langle 1 \rangle 2.$ $\bigcup \mathcal{P} \mathbf{A} \subseteq \mathbf{A}$
   $\langle 2 \rangle 1.$ LET: $x \in \bigcup \mathcal{P} \mathbf{A}$
   $\langle 2 \rangle 2.$ PICK $X \in \mathcal{P} \mathbf{A}$ such that $x \in X$
   $\langle 2 \rangle 3.$ $X \subseteq \mathbf{A}$
   $\langle 2 \rangle 4.$ $x \in \mathbf{A}$
☐

1.6   Transitive Classes

Definition 1.6.1 (Transitive Class)**.** A class $\mathbf{A}$ is a *transitive class* iff whenever $x \in y \in \mathbf{A}$ then $x \in \mathbf{A}$.

Proposition Schema 1.6.2 (FOL)**.** *For any class* $\mathbf{A}$*, the following is a theorem:*
   The following are equivalent.

   1. $\mathbf{A}$ *is a transitive class.*

   2. $\bigcup \mathbf{A} \subseteq \mathbf{A}$

   3. *Every element of* $\mathbf{A}$ *is a subset of* $\mathbf{A}$*.*

   4. $\mathbf{A} \subseteq \mathcal{P} \mathbf{A}$

PROOF: Immediate from definitions. ☐

Proposition Schema 1.6.3 (FOL)**.** *For any class* $\mathbf{A}$*, the following is a theorem:*
   If $\mathbf{A}$ *is a transitive class then* $\bigcup \mathbf{A}$ *is a transitive class.*

PROOF:
$\langle 1 \rangle 1.$ ASSUME: $\mathbf{A}$ is a transitive class.
$\langle 1 \rangle 2.$ LET: $x \in y \in \bigcup \mathbf{A}$
$\langle 1 \rangle 3.$ $y \in \mathbf{A}$
   PROOF: Since $\bigcup \mathbf{A} \subseteq \mathbf{A}$ by Proposition 1.6.2.
$\langle 1 \rangle 4.$ $x \in \bigcup \mathbf{A}$
☐

Proposition Schema 1.6.4 (Z)**.** *For any class* $\mathbf{A}$*, the following is a theorem:*
   We have $\mathbf{A}$ *is a transitive class if and only if* $\mathcal{P} \mathbf{A}$ *is a transitive class.*

PROOF:
$\langle 1 \rangle 1.$ If $\mathbf{A}$ is a transitive class then $\mathcal{P} \mathbf{A}$ is a transitive class.
   $\langle 2 \rangle 1.$ ASSUME: $\mathbf{A}$ is a transitive class.
   $\langle 2 \rangle 2.$ $\mathbf{A} \subseteq \mathcal{P} \mathbf{A}$
      PROOF: Proposition 1.6.2.
   $\langle 2 \rangle 3.$ $\mathcal{P} \mathbf{A} \subseteq \mathcal{P} \mathcal{P} \mathbf{A}$

PROOF: Proposition 1.5.8.
⟨2⟩4. $\mathcal{P}\mathbf{A}$ is a transitive class.
PROOF: Proposition 1.6.2.
⟨1⟩2. If $\mathcal{P}\mathbf{A}$ is a transitive class then $\mathbf{A}$ is a transitive class.
⟨2⟩1. ASSUME: $\mathcal{P}\mathbf{A}$ is a transitive class.
⟨2⟩2. $\bigcup\mathcal{P}\mathbf{A} \subseteq \mathcal{P}\mathbf{A}$
PROOF: Proposition 1.6.2.
⟨2⟩3. $\mathbf{A} \subseteq \mathcal{P}\mathbf{A}$
PROOF: Proposition 1.5.10.
⟨2⟩4. $\mathbf{A}$ is a transitive class.
PROOF: Proposition 1.6.2.
□

Proposition Schema 1.6.5 (FOL)**.** *For any class* $\mathbf{A}$*, the following is a theorem:*

If every member of $\mathbf{A}$ *is a transitive set then* $\bigcup\mathbf{A}$ *is a transitive class.*

PROOF:
⟨1⟩1. ASSUME: Every member of $\mathbf{A}$ is a transitive set.
⟨1⟩2. LET: $x \in y \in \bigcup\mathbf{A}$
⟨1⟩3. PICK $A \in \mathbf{A}$ such that $y \in A$.
⟨1⟩4. $x \in A$
PROOF: Since A is a transitive set.
⟨1⟩5. $x \in \bigcup\mathbf{A}$
□

Proposition Schema 1.6.6 (FOL)**.** *For any class* $\mathbf{A}$*, the following is a theorem:*

If every member of $\mathbf{A}$ *is a transitive set then* $\bigcap\mathbf{A}$ *is a transitive class.*

PROOF:
⟨1⟩1. ASSUME: Every member of $\mathbf{A}$ is a transitive set.
⟨1⟩2. LET: $x \in y \in \bigcap\mathbf{A}$
PROVE: $x \in \bigcap\mathbf{A}$
⟨1⟩3. LET: $A \in \mathbf{A}$
⟨1⟩4. $y \in A$
⟨1⟩5. $x \in A$
PROOF: Since A is a transitive set.
□

Chapter 2

Relations

2.1 Ordered Pairs

Definition 2.1.1 (Ordered Pair)**.** For any sets a and b, the *ordered pair* (a, b) is defined to be $\{\{a\}, \{a, b\}\}$.

Theorem 2.1.2 (Z)**.** *For any sets a, b, c, d, we have $(a, b) = (c, d)$ if and only if $a = c$ and $b = d$.*

PROOF:
$\langle 1 \rangle 1$. If $(a, b) = (c, d)$ then $a = c$ and $b = d$.
   $\langle 2 \rangle 1$. ASSUME: $\{\{a\}, \{a, b\}\} = \{\{c\}, \{c, d\}\}$
   $\langle 2 \rangle 2$. $\bigcap\{\{a\}, \{a, b\}\} = \bigcap\{\{c\}, \{c, d\}\}$
   $\langle 2 \rangle 3$. $\{a\} = \{c\}$
   $\langle 2 \rangle 4$. $a = c$
   $\langle 2 \rangle 5$. $\bigcup\{\{a\}, \{a, b\}\} = \bigcup\{\{c\}, \{c, d\}\}$
   $\langle 2 \rangle 6$. $\{a, b\} = \{c, d\}$
   $\langle 2 \rangle 7$. $b = c$ or $b = d$
   $\langle 2 \rangle 8$. $a = d$ or $b = d$
   $\langle 2 \rangle 9$. If $b = c$ and $a = d$ then $b = d$
     PROOF: By $\langle 2 \rangle 4$.
   $\langle 2 \rangle 10$. $b = d$
     PROOF: From $\langle 2 \rangle 7$, $\langle 2 \rangle 8$, $\langle 2 \rangle 9$.
$\langle 1 \rangle 2$. If $a = c$ and $b = d$ then $(a, b) = (c, d)$.
   PROOF: First-order logic.
$\square$

Definition 2.1.3 (Cartesian Product)**.** The *Cartesian product* of classes $\mathbf{A}$ and $\mathbf{B}$ is the class $\mathbf{A} \times \mathbf{B} := \{(x, y) \mid x \in \mathbf{A}, y \in \mathbf{B}\}$.

Proposition 2.1.4 (Z)**.** *For any sets A and B, the class $A \times B$ is a set.*

PROOF: It is a subset of $\mathcal{P}\mathcal{P}(A \cup B)$. $\square$

Proposition Schema 2.1.5 (Z)**.** *For any classes* $\mathbf{A}$, $\mathbf{B}$ *and* $\mathbf{C}$, *the following is a theorem:*

$$\mathbf{A} \times (\mathbf{B} \cup \mathbf{C}) = (\mathbf{A} \times \mathbf{B}) \cup (\mathbf{A} \times \mathbf{C})$$

PROOF:
$$
\begin{aligned}
(x,y) \in \mathbf{A} \times (\mathbf{B} \cup \mathbf{C}) &\Leftrightarrow x \in \mathbf{A} \wedge (y \in \mathbf{B} \vee y \in \mathbf{C}) \\
&\Leftrightarrow (x \in \mathbf{A} \wedge y \in \mathbf{B}) \vee (x \in \mathbf{A} \wedge y \in \mathbf{C}) \\
&\Leftrightarrow (x,y) \in (\mathbf{A} \times \mathbf{B}) \cup (\mathbf{A} \times \mathbf{C}) \qquad \Box
\end{aligned}
$$

Proposition Schema 2.1.6 (Z)**.** *For any classes* $\mathbf{A}$ *and* $\mathbf{B}$, *the following is a theorem:*

If $\mathbf{A} \times \mathbf{B} = \mathbf{A} \times \mathbf{C}$ *and* $\mathbf{A}$ *is nonempty then* $\mathbf{B} = \mathbf{C}$.

PROOF:
$\langle 1 \rangle 1.$ PICK $a \in \mathbf{A}$
$\langle 1 \rangle 2.$ For all x we have $x \in \mathbf{B}$ iff $x \in \mathbf{C}$.
   PROOF:
$$
\begin{aligned}
x \in \mathbf{B} &\Leftrightarrow (a,x) \in \mathbf{A} \times \mathbf{B} \\
&\Leftrightarrow (a,x) \in \mathbf{A} \times \mathbf{C} \\
&\Leftrightarrow x \in \mathbf{C}
\end{aligned}
$$

$\Box$

Proposition Schema 2.1.7 (Z)**.** *For any classes* $\mathbf{A}$ *and* $\mathbf{B}$, *the following is a theorem:*

$$\mathbf{A} \times \bigcup \mathbf{B} = \{(a,b) \mid \exists Y \in \mathbf{B}.(a \in \mathbf{A} \wedge b \in Y)\}$$

PROOF:
$$
\begin{aligned}
(x,y) \in A \times \bigcup \mathbf{B} &\Leftrightarrow x \in A \wedge \exists Y \in \mathbf{B}.y \in Y \\
&\Leftrightarrow \exists Y \in \mathbf{B}(x \in A \wedge y \in Y) \qquad \Box
\end{aligned}
$$

2.2   Relations

Definition 2.2.1 (Relation)**.** A *relation* $\mathbf{R}$ between classes $\mathbf{A}$ and $\mathbf{B}$ is a subclass of $\mathbf{A} \times \mathbf{B}$.

A *(binary) relation on* $\mathbf{A}$ is a relation between $\mathbf{A}$ and $\mathbf{A}$.

We write $x\mathbf{R}y$ for $(x,y) \in \mathbf{R}$.

2.2.1   Identity Functions

Definition 2.2.2 (Identity Function)**.** For any class $\mathbf{A}$, the *identity function* or *diagonal relation* $\mathrm{id}_{\mathbf{A}}$ on $\mathbf{A}$ is

$$\mathrm{id}_{\mathbf{A}} := \{(x,x) \mid x \in \mathbf{A}\} \ .$$

2.2.2 Inverses

Definition 2.2.3 (Inverse)**.** The *inverse* of a relation $\mathbf{R}$ between $\mathbf{A}$ and $\mathbf{B}$ is the relation $\mathbf{R}^{-1}$ between $\mathbf{B}$ and $\mathbf{A}$ defined by

$$b\mathbf{R}^{-1}a \Leftrightarrow a\mathbf{R}b \ .$$

Proposition Schema 2.2.4 (Z)**.** *For any classes* $\mathbf{A}$*,* $\mathbf{B}$ *and* $\mathbf{R}$*, the following is a theorem:*

   If $\mathbf{R}$ *is a relation between* $\mathbf{A}$ *and* $\mathbf{B}$*, we have* $(\mathbf{R}^{-1})^{-1} = \mathbf{R}$*.*

PROOF:
$$x(\mathbf{R}^{-1})^{-1}y \Leftrightarrow y\mathbf{R}^{-1}x$$
$$\Leftrightarrow x\mathbf{R}y \qquad\qquad \Box$$

2.2.3 Composition

Definition 2.2.5 (Composition)**.** Let $\mathbf{R}$ be a relation between $\mathbf{A}$ and $\mathbf{B}$, and $\mathbf{S}$ be a relation between $\mathbf{B}$ and $\mathbf{C}$. The *composition* $\mathbf{S} \circ \mathbf{R}$ is the relation between $\mathbf{A}$ and $\mathbf{C}$ defined by
$$a(\mathbf{S} \circ \mathbf{R})c \Leftrightarrow \exists b(a\mathbf{R}b \wedge b\mathbf{S}c) \ .$$

Proposition Schema 2.2.6 (Z)**.** *For any classes* $\mathbf{A}$*,* $\mathbf{B}$*,* $\mathbf{C}$*,* $\mathbf{R}$ *and* $\mathbf{S}$*, the following is a theorem:*

   If $\mathbf{R}$ *is a relation between* $\mathbf{A}$ *and* $\mathbf{B}$*, and* $\mathbf{S}$ *is a relation between* $\mathbf{B}$ *and* $\mathbf{C}$*, then*
$$(\mathbf{S} \circ \mathbf{R})^{-1} = \mathbf{R}^{-1} \circ \mathbf{S}^{-1} \ .$$

PROOF:
$$z(\mathbf{S} \circ \mathbf{R})^{-1}x \Leftrightarrow x(\mathbf{S} \circ \mathbf{R})z$$
$$\Leftrightarrow \exists y.(x\mathbf{R}y \wedge y\mathbf{S}z)$$
$$\Leftrightarrow \exists y.(y\mathbf{R}^{-1}x \wedge z\mathbf{S}^{-1}y)$$
$$\Leftrightarrow z(\mathbf{R}^{-1} \circ \mathbf{S}^{-1})x \qquad\qquad \Box$$

2.2.4 Properties of Relaitons

Definition 2.2.7 (Reflexive)**.** Let $\mathbf{R}$ be a binary relation on $\mathbf{A}$. Then $\mathbf{R}$ is *reflexive* on $\mathbf{A}$ iff $\forall x \in \mathbf{A}.(x,x) \in \mathbf{R}$.

Proposition Schema 2.2.8 (Z)**.** *For any classes* $\mathbf{A}$ *and* $\mathbf{R}$*, the following is a theorem:*

   If $\mathbf{R}$ *is a reflexive relation on* $\mathbf{A}$ *then so is* $\mathbf{R}^{-1}$*.*

PROOF:
⟨1⟩1. LET: $x \in \mathbf{A}$
⟨1⟩2. $x\mathbf{R}x$
   PROOF: Since $\mathbf{R}$ is reflexive.

$\langle 1 \rangle 3.$ $x\mathbf{R}^{-1}x$
☐

Definition 2.2.9 (Irreflexive)**.** A relation $\mathbf{R}$ is *irreflexive* iff there is no x such that $(x, x) \in \mathbf{R}$.

Definition 2.2.10 (Symmetric)**.** A relation $\mathbf{R}$ is *symmetric* iff, whenever $x\mathbf{R}y$, then $y\mathbf{R}x$.

Definition 2.2.11 (Antisymmetric)**.** A relation $\mathbf{R}$ is *antisymmetric* iff, whenever $x\mathbf{R}y$ and $y\mathbf{R}x$, then $x = y$.

Proposition Schema 2.2.12 (Z)**.** *For any classes* $\mathbf{A}$ *and* $\mathbf{R}$*, the following is a theorem:*

If $\mathbf{R}$ *is an antisymmetric relation on* $\mathbf{A}$ *then so is* $\mathbf{R}^{-1}$*.*

PROOF:
$\langle 1 \rangle 1.$ ASSUME: $x\mathbf{R}^{-1}y$ and $y\mathbf{R}^{-1}x$
$\langle 1 \rangle 2.$ $y\mathbf{R}x$ and $x\mathbf{R}y$
$\langle 1 \rangle 3.$ $x = y$
   PROOF: Since $\mathbf{R}$ is antisymmetric.
☐

Definition 2.2.13 (Transitive)**.** A relation $\mathbf{R}$ is *transitive* iff, whenever $x\mathbf{R}y$ and $y\mathbf{R}z$, then $x\mathbf{R}z$.

Proposition Schema 2.2.14 (Z)**.** *For any classes* $\mathbf{A}$*,* $\mathbf{B}$ *and* $\mathbf{R}$*, the following is a theorem:*

If $\mathbf{R}$ *is a transitive relation between* $\mathbf{A}$ *and* $\mathbf{B}$ *then* $\mathbf{R}^{-1}$ *is transitive.*

PROOF:
$\langle 1 \rangle 1.$ ASSUME: $(x, y), (y, z) \in \mathbf{R}^{-1}$
$\langle 1 \rangle 2.$ $(y, x), (z, y) \in \mathbf{R}$
$\langle 1 \rangle 3.$ $(z, x) \in \mathbf{R}$
$\langle 1 \rangle 4.$ $(x, z) \in \mathbf{R}^{-1}$
☐

Proposition 2.2.15 (Z)**.** *For any relation* R *on a set* A*, there exists a smallest transitive relation on* A *that includes* R*.*

PROOF: The relation is $\bigcap \{S \in \mathcal{P}A^2 \mid R \subseteq S, S \text{ is transitive}\}$. ☐

Definition 2.2.16 (Transitive Closure)**.** For any relation R on a set A, the *transitive closure* of R is the smallest transitive relation that includes R.

Definition 2.2.17 (Minimal)**.** Let $\mathbf{R}$ be a relation on $\mathbf{A}$. An element $m \in \mathbf{A}$ is *minimal* iff there is no $x \in \mathbf{A}$ such that $x\mathbf{R}m$.

Definition 2.2.18 (Maximal)**.** Let $\mathbf{R}$ be a relation on $\mathbf{A}$. An element $m \in \mathbf{A}$ is *maximal* iff there is no $x \in \mathbf{A}$ such that $m\mathbf{R}x$.

2.3   n-ary Relations

Definition Schema 2.3.1. For any sets $a_1, \ldots, a_n$, define the *ordered n-tuple* $(a_1, \ldots, a_n)$ by

$$(a_1) := a_1$$
$$(a_1, \ldots, a_n, a_{n+1}) = ((a_1, \ldots, a_n), a_{n+1})$$

Definition Schema 2.3.2. An *n-ary relation on* $\mathbf{A}$ is a class of ordered n-tuples all of whose components are in $\mathbf{A}$.

2.4   Well Founded Relations

Definition 2.4.1 (Well Founded)**.** A relation $\mathbf{R}$ on a class $\mathbf{A}$ is *well founded* iff:

- for all $a \in A$, the class $\{x \in \mathbf{A} \mid x\mathbf{R}a\}$ is a set;

- every nonempty subset of $\mathbf{A}$ has an $\mathbf{R}$-minimal element.

Proposition 2.4.2 (Z)**.** *For any class* $\mathbf{A}$*, the relation* $\{(x, y) \in \mathbf{A}^2 \mid x \in y\}$ *is well founded.*

PROOF:
$\langle 1 \rangle 1$. For all $a \in \mathbf{A}$, the class $\{x \in \mathbf{A} \mid x \in a\}$ is a set.
  PROOF: It is a subclass of a.
$\langle 1 \rangle 2$. Every nonempty subset of $\mathbf{A}$ has an $\in$-minimal element.
  $\langle 2 \rangle 1$. LET: C be a nonempty subset of $\mathbf{A}$
  $\langle 2 \rangle 2$. PICK $m \in C$ such that $m \cap C = \emptyset$
    PROOF: Axiom of Regularity.
  $\langle 2 \rangle 3$. m is $\in$-minimal in C.
☐

Proposition Schema 2.4.3 (Z)**.** *For any classes* $\mathbf{A}$*,* $\mathbf{B}$ *and* $\mathbf{R}$*, the following is a theorem:*
  Assume $\mathbf{R}$ *is a well founded relation on* $\mathbf{A}$ *and* $\mathbf{B} \subseteq \mathbf{A}$ *is nonempty. Then* $\mathbf{B}$ *has an* $\mathbf{R}$*-minimal element.*

PROOF:
$\langle 1 \rangle 1$. PICK $b \in \mathbf{B}$
$\langle 1 \rangle 2$. LET: $S = \{x \in \mathbf{B} \mid x\mathbf{R}b\}$
  PROOF: S is a set because it is a subclass of $\{x \in \mathbf{A} \mid x\mathbf{R}b\}$.
$\langle 1 \rangle 3$. CASE: $S = \emptyset$
  PROOF: In this case b is an $\mathbf{R}$-minimal element of $\mathbf{B}$.
$\langle 1 \rangle 4$. CASE: $S \neq \emptyset$
  PROOF: In this cases S has an $\mathbf{R}$-minimal element, which is an $\mathbf{R}$-minimal element of $\mathbf{B}$.
☐

Proposition Schema 2.4.4 (Z)**.** *For any classes* **A***,* **B** *and* **R***, the following is a theorem:*

   Assume **R** *is a well founded relation on* **B** *and* **A** $\subseteq$ **B***. Then* **R** $\cap$ **A**^{2} *is a well founded relation on* **A***.*

PROOF:

$\langle 1 \rangle$1. LET: $\mathbf{R}' = \mathbf{R} \cap \mathbf{A}^2$

$\langle 1 \rangle$2. For all $a \in \mathbf{A}$, the class $\{x \in \mathbf{A} \mid x\mathbf{R}'a\}$ is a set.

   PROOF: By Comprehension since it is a subclass of $\{x \in \mathbf{B} \mid x\mathbf{R}a\}$.

$\langle 1 \rangle$3. Every nonempty subset of **A** has an **R**$'$-minimal element.

   PROOF: It is a nonempty subset of **B** and so has an **R**-minimal element, which is also an **R**$'$-minimal element.

$\square$

Theorem Schema 2.4.5 (Transfinite Induction Principle (Z))**.** *For any classes* **A***,* **B** *and* **R***, the following is a theorem:*

   Assume **R** *is a well founded relation on* **A** *and* **B** $\subseteq$ **A***. Assume that, for all* $t \in \mathbf{A}$,

$$\{x \in \mathbf{A} \mid x\mathbf{R}t\} \subseteq \mathbf{B} \Rightarrow t \in \mathbf{B} \ .$$

Then **B** = **A***.*

PROOF:

$\langle 1 \rangle$1. ASSUME: for a contradiction **B** $\neq$ **A**

$\langle 1 \rangle$2. PICK an **R**-minimal element m of **A** − **B**.

   PROOF: Proposition 2.4.3.

$\langle 1 \rangle$3. $\{x \in \mathbf{A} \mid x\mathbf{R}m\} \subseteq \mathbf{B}$

   PROOF: By minimality of m.

$\langle 1 \rangle$4. $m \in \mathbf{B}$

$\langle 1 \rangle$5. Q.E.D.

   PROOF: This is a contradiction.

$\square$

Theorem 2.4.6 (Z)**.** *The transitive closure of a well founded relation on a set is well founded.*

PROOF:

$\langle 1 \rangle$1. LET: R be a well founded relation on the set A.

$\langle 1 \rangle$2. LET: R^t be the transitive closure of R.

$\langle 1 \rangle$3. For any $x, y \in A$, if xR^ty then there exists $z \in A$ such that zRy.

   PROOF: $\{(x, y) \in A^2 \mid \exists z \in A. zRy\}$ is a transitive relation on A that includes R.

$\langle 1 \rangle$4. LET: B be a nonempty subset of A.

$\langle 1 \rangle$5. PICK an R-minimal element b of B.

$\langle 1 \rangle$6. b is R^t-minimal in B.

   PROOF: If there exists x such that xR^tb then there exists z such that zRb by $\langle 1 \rangle$3.

$\square$

Definition 2.4.7 (Initial Segment)**.** Let $\mathbf{R}$ be a relation on $\mathbf{A}$ and $a \in \mathbf{A}$. The *initial segment* up to a is

$$\operatorname{seg} a := \{x \in \mathbf{A} \mid x\mathbf{R}a\} \ .$$

Theorem Schema 2.4.8 (Transfinite Recursion Theorem Schema (ZFC))**.** *For any classes $\mathbf{A}$, $\mathbf{R}$ and any property $G[x, y, z]$, there exists a class $\mathbf{F}$ such that, for any class $\mathbf{F}'$ the following is a theorem:*

Assume that $\mathbf{R}$ is a well-founded relation on $\mathbf{A}$. Assume that, for any f and t, there exists a unique z such that $G[f, t, z]$. Then $\mathbf{F} : \mathbf{A} \to \mathbf{V}$ such that, for all $t \in \mathbf{A}$, we have $\mathbf{F} \restriction \operatorname{seg} t$ is a set and

$$G[\mathbf{F} \restriction \operatorname{seg} t, t, \mathbf{F}(t)] \ .$$

If $\mathbf{F}' : \mathbf{A} \to \mathbf{V}$ satisfies that, for all $t \in \mathbf{A}$, we have $\mathbf{F}' \restriction \operatorname{seg} t$ is a set and $G[\mathbf{F}' \restriction \operatorname{seg} t, t, \mathbf{F}'(t)]$, then $\mathbf{F}' = \mathbf{F}$.

PROOF:

$\langle 1 \rangle 1$. For B a subset of $\mathbf{A}$, let us say a function $v : B \to \mathbf{V}$ is *acceptable* iff, for all $x \in B$, we have $\operatorname{seg} x \subseteq B$ and $G[v \restriction \operatorname{seg} x, x, v(x)]$

$\langle 1 \rangle 2$. LET: $\mathbf{K}$ be the class of all acceptable functions.

$\langle 1 \rangle 3$. LET: $\mathbf{F} = \bigcup \mathbf{K}$

$\langle 1 \rangle 4$. For all $B, C \subseteq \mathbf{A}$, given $v_1 : B \to \mathbf{V}$ and $v_2 : C \to \mathbf{V}$ acceptable and $x \in B \cap C$, we have $v_1(x) = v_2(x)$

$\quad \langle 2 \rangle 1$. ASSUME: as transfinite induction hypothesis $\forall y \mathbf{R} x.y \in B \cap C \Rightarrow v_1(y) = v_2(y)$

$\quad \langle 2 \rangle 2$. $v_1 \restriction \operatorname{seg} x = v_2 \restriction \operatorname{seg} x$

$\quad \langle 2 \rangle 3$. $G[v_1 \restriction \operatorname{seg} x, x, v_1(x)]$

$\quad \langle 2 \rangle 4$. $G[v_2 \restriction \operatorname{seg} x, x, v_2(x)]$

$\quad \langle 2 \rangle 5$. $v_1(x) = v_2(x)$

$\langle 1 \rangle 5$. $\mathbf{F}$ is a function.

$\quad \langle 2 \rangle 1$. ASSUME: $(x, y), (x, z) \in \mathbf{F}$

$\quad \langle 2 \rangle 2$. PICK acceptable $v_1 : B \to \mathbf{V}$ and $v_2 : C \to \mathbf{V}$ such that $v_1(x) = y$ and $v_2(x) = z$

$\quad \langle 2 \rangle 3$. $y = z$

$\quad\quad$ PROOF: By $\langle 1 \rangle 4$.

$\langle 1 \rangle 6$. For all $t \in \operatorname{dom} \mathbf{F}$, we have $\mathbf{F} \restriction \operatorname{seg} t$ is a set and $G[\mathbf{F} \restriction \operatorname{seg} t, t, \mathbf{F}(t)]$

$\quad \langle 2 \rangle 1$. LET: $t \in \operatorname{dom} \mathbf{F}$

$\quad \langle 2 \rangle 2$. PICK an acceptable $v : A \to \mathbf{V}$ such that $t \in A$

$\quad \langle 2 \rangle 3$. For all $y \mathbf{R} x$ we have $v(y) = \mathbf{F}(y)$

$\quad \langle 2 \rangle 4$. $\mathbf{F} \restriction \operatorname{seg} x = v \restriction \operatorname{seg} x$

$\quad \langle 2 \rangle 5$. $G[v \restriction \operatorname{seg} x, x, v(x)]$

$\quad \langle 2 \rangle 6$. $G[\mathbf{F} \restriction \operatorname{seg} x, x, \mathbf{F}(x)]$

$\langle 1 \rangle 7$. $\operatorname{dom} \mathbf{F} = \mathbf{A}$

$\quad \langle 2 \rangle 1$. LET: $x \in \mathbf{A}$

$\quad \langle 2 \rangle 2$. ASSUME: as transfinite induction hypothesis $\forall y \mathbf{R} x.y \in \mathbf{A}$

$\quad \langle 2 \rangle 3$. ASSUME: for a contradiction $x \notin \operatorname{dom} \mathbf{F}$

⟨2⟩4. $\mathbf{F} \upharpoonright \operatorname{seg} x$ is a set
    Proof: Axiom of Replacement.
⟨2⟩5. $\mathbf{F} \upharpoonright \operatorname{seg} x$ is acceptable
⟨2⟩6. Let: y be the unique object such that $G[\mathbf{F} \upharpoonright \operatorname{seg} x, x, y]$
⟨2⟩7. $\mathbf{F} \upharpoonright \operatorname{seg} x \cup \{(x, y)\}$ is acceptable
⟨2⟩8. $x \in \operatorname{dom} \mathbf{F}$
⟨2⟩9. Q.E.D.
    Proof: This is a contradiction.
⟨1⟩8. If $\mathbf{F}' : \mathbf{A} \to \mathbf{V}$ satisfies the theorem, then $\mathbf{F}' = \mathbf{F}$.
  ⟨2⟩1. Let: $x \in \mathbf{A}$
      Prove: $\mathbf{F}'(x) = \mathbf{F}(x)$
  ⟨2⟩2. Assume: as transfinite induction hypothesis $\forall y \mathbf{R} x.\mathbf{F}'(y) = \mathbf{F}(y)$
  ⟨2⟩3. $\mathbf{F} \upharpoonright x = \mathbf{F}' \upharpoonright x$
  ⟨2⟩4. $G[\mathbf{F} \upharpoonright x, x, \mathbf{F}(x)]$
  ⟨2⟩5. $G[\mathbf{F}' \upharpoonright x, x, \mathbf{F}'(x)]$
  ⟨2⟩6. $\mathbf{F}(x) = \mathbf{F}'(x)$
▯

Chapter 3

Functions

3.1 Functions

Definition 3.1.1 (Function). A *function* from $\mathbf{A}$ to $\mathbf{B}$ is a relation $\mathbf{F}$ between $\mathbf{A}$ and $\mathbf{B}$ such that, for all $x \in \mathbf{A}$, there is only one y such that $x\mathbf{F}y$. We denote this y by $\mathbf{F}(x)$.

A *binary operation* on a class $\mathbf{A}$ is a function $\mathbf{A}^2 \to \mathbf{A}$.

Definition 3.1.2 (Closed). Let $\mathbf{F} : \mathbf{A} \to \mathbf{A}$ be a function and $\mathbf{B} \subseteq \mathbf{A}$. Then $\mathbf{B}$ is *closed* under $\mathbf{F}$ iff $\forall x \in \mathbf{B}.\mathbf{F}(x) \in \mathbf{B}$.

Proposition 3.1.3 (Z). *For any class $\mathbf{A}$, the following is a theorem:*

$$\mathrm{id}_{\mathbf{A}} : \mathbf{A} \to \mathbf{A}$$

PROOF: For all $x \in \mathbf{A}$, the only y such that $(x, y) \in \mathrm{id}_{\mathbf{A}}$ is $y = x$. $\square$

Proposition Schema 3.1.4 (Z). *For any classes $\mathbf{A}$, $\mathbf{B}$, $\mathbf{C}$, $\mathbf{F}$ and $\mathbf{G}$, the following is a theorem:*

Assume $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ and $\mathbf{G} : \mathbf{B} \to \mathbf{C}$. Then $\mathbf{G} \circ \mathbf{F} : \mathbf{A} \to \mathbf{C}$ and, for all $x \in \mathbf{A}$, we have

$$(\mathbf{G} \circ \mathbf{F})(x) = \mathbf{G}(\mathbf{F}(x)) \ .$$

PROOF:
⟨1⟩1. $\forall x \in \mathbf{A}.(x, \mathbf{G}(\mathbf{F}(x))) \in \mathbf{G} \circ \mathbf{F}$
   PROOF: Because $(x, \mathbf{F}(x)) \in \mathbf{F}$ and $(\mathbf{F}(x), \mathbf{G}(\mathbf{F}(x))) \in \mathbf{G}$.
⟨1⟩2. If $(x, z) \in \mathbf{F} \circ \mathbf{G}$ then $z = \mathbf{G}(\mathbf{F}(x))$
   ⟨2⟩1. PICK $y \in \mathbf{B}$ such that $x\mathbf{F}y$ and $y\mathbf{G}z$
   ⟨2⟩2. $y = \mathbf{F}(x)$
   ⟨2⟩3. $z = \mathbf{G}(y)$
   ⟨2⟩4. $z = \mathbf{G}(\mathbf{F}(x))$
$\square$

Proposition 3.1.5 (Z)**.** *For any set A there exists a function $F : \mathcal{P}A - \{\emptyset\} \to A$ (a* choice function *for A) such that, for every nonempty $B \subseteq A$, we have $F(B) \in B$.*

PROOF:

$\langle 1 \rangle 1.$ LET: A be a set.

$\langle 1 \rangle 2.$ LET: $\mathcal{A} = \{\{B\} \times B \mid B \in \mathcal{P}A - \{\emptyset\}\}$

$\langle 1 \rangle 3.$ Every member of $\mathcal{A}$ is nonempty.

$\langle 1 \rangle 4.$ Any two distinct members of $\mathcal{A}$ are disjoint.

$\langle 1 \rangle 5.$ PICK a set C such that, for all $X \in \mathcal{A}$, we have $C \cap X$ is a singleton.

   PROOF: Axiom of Choice.

$\langle 1 \rangle 6.$ LET: $F = C \cap \bigcup \mathcal{A}$

$\langle 1 \rangle 7.$ $F : \mathcal{P}A - \{\emptyset\} \to A$

   $\langle 2 \rangle 1.$ F is a function.

      $\langle 3 \rangle 1.$ LET: $(B, b), (B, b') \in F$

      $\langle 3 \rangle 2.$ $(B, b), (B, b') \in \{B\} \times B$

         PROOF: Since $(B, b), (B, b') \in \bigcup \mathcal{A}$.

      $\langle 3 \rangle 3.$ $(B, b), (B, b') \in C \cap (\{B\} \times B)$

      $\langle 3 \rangle 4.$ $(B, b) = (B, b')$

         PROOF: From $\langle 1 \rangle 5$.

      $\langle 3 \rangle 5.$ $b = b'$

   $\langle 2 \rangle 2.$ $\operatorname{dom} F = \mathcal{P}A - \{\emptyset\}$

      PROOF:

$$
\begin{aligned}
& B \in \operatorname{dom} F \\
\Leftrightarrow & \exists b.(B, b) \in F \\
\Leftrightarrow & \exists b.((B, b) \in \bigcup \mathcal{A} \wedge (B, b) \in C) \\
\Leftrightarrow & \exists b.\exists B' \in \mathcal{P}A - \{\emptyset\}.((B, b) \in \{B'\} \times B' \wedge (B, b) \in C) \\
\Leftrightarrow & B \in \mathcal{P}A - \{\emptyset\} \wedge \exists b \in B.(B, b) \in C \\
\Leftrightarrow & B \in \mathcal{P}A - \{\emptyset\} \qquad\qquad\qquad\qquad\qquad\qquad (\langle 1 \rangle 5)
\end{aligned}
$$

   $\langle 2 \rangle 3.$ $\operatorname{ran} F \subseteq A$

$\langle 1 \rangle 8.$ For every nonempty $B \subseteq A$ we have $F(B) \in B$

▯

Proposition 3.1.6 (Z)**.** *For any relation R between A and B, there exists a function $H : A \to B$ such that $H \subseteq R$ (i.e. $\forall x \in A.xRH(x)$).*

PROOF:

$\langle 1 \rangle 1.$ LET: R be a relation between A and B.

$\langle 1 \rangle 2.$ PICK a choice function G for B.

$\langle 1 \rangle 3.$ Define $H : A \to B$ by $H(x) = G(\{y \mid xRy\})$

$\langle 1 \rangle 4.$ $H \subseteq R$

▯

3.1.1 Injective Functions

Definition 3.1.7 (Injective)**.** A function $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ is *one-to-one*, *injective* or an *injection*, $\mathbf{F} : \mathbf{A} \rightarrowtail \mathbf{B}$, iff, for all $x, y \in \mathbf{A}$, if $\mathbf{F}(x) = \mathbf{F}(y)$, then $x = y$.

Proposition 3.1.8 (Z)**.** *For any class* $\mathbf{A}$, *the following is a theorem:*
   $\mathrm{id}_{\mathbf{A}} : \mathbf{A} \to \mathbf{A}$ *is injective.*

PROOF: If $\mathrm{id}_{\mathbf{A}}(x) = \mathrm{id}_{\mathbf{A}}(y)$ then immediately $x = y$. □

Proposition Schema 3.1.9 (Z)**.** *For any classes* $\mathbf{A}$, $\mathbf{B}$, $\mathbf{C}$, $\mathbf{F}$, $\mathbf{G}$, *the following is a theorem:*
   Assume $\mathbf{F} : \mathbf{A} \rightarrowtail \mathbf{B}$ *and* $\mathbf{G} : \mathbf{B} \rightarrowtail \mathbf{C}$. *Then* $\mathbf{G} \circ \mathbf{F} : \mathbf{A} \rightarrowtail \mathbf{C}$.

PROOF:
$\langle 1 \rangle 1$. LET: $x, y \in \mathbf{A}$
$\langle 1 \rangle 2$. ASSUME: $(\mathbf{G} \circ \mathbf{F})(x) = (\mathbf{G} \circ \mathbf{F})(y)$
$\langle 1 \rangle 3$. $\mathbf{G}(\mathbf{F}(x)) = \mathbf{G}(\mathbf{F}(y))$
$\langle 1 \rangle 4$. $\mathbf{F}(x) = \mathbf{F}(y)$
   PROOF: Since $\mathbf{G}$ is injective.
$\langle 1 \rangle 5$. $x = y$
   PROOF: Since $\mathbf{F}$ is injective.
□

Proposition 3.1.10 (Z)**.** *Let* $F : A \to B$ *where* A *is nonempty. There exists* $G : B \to A$ *(a* left inverse*) such that* $G \circ F = \mathrm{id}_A$ *if and only if* F *is one-to-one.*

PROOF:
$\langle 1 \rangle 1$. If there exists $G : B \to A$ such that $G \circ F = \mathrm{id}_A$ then F is one-to-one.
   $\langle 2 \rangle 1$. ASSUME: $G : B \to A$ and $G \circ F = I_A$
   $\langle 2 \rangle 2$. LET: $x, y \in A$
   $\langle 2 \rangle 3$. ASSUME: $F(x) = F(y)$
   $\langle 2 \rangle 4$. $x = y$
      PROOF: $x = G(F(x)) = G(F(y)) = y$
$\langle 1 \rangle 2$. If F is one-to-one then there exists $G : B \to A$ such that $G \circ F = I_A$.
   $\langle 2 \rangle 1$. ASSUME: F is one-to-one.
   $\langle 2 \rangle 2$. PICK $a \in A$
   $\langle 2 \rangle 3$. LET: $G : B \to A$ be the function defined by: $G(b)$ is the (unique) $x \in A$
            such that $F(x) = b$ if there exists such an x, $G(b) = a$ otherwise.
   $\langle 2 \rangle 4$. For all $x \in A$ we have $G(F(x)) = x$.
□

3.1.2 Surjective Functions

Definition 3.1.11 (Surjective)**.** Let $F : A \to B$. We say that F is *surjective*, or maps A *onto* B, and write $F : A \twoheadrightarrow B$, iff for all $y \in B$ there exists $x \in A$ such that $F(x) = y$.

Proposition Schema 3.1.12 (Z)**.** *For any class* $\mathbf{A}$, *the following is a theorem:*
   $\mathrm{id}_{\mathbf{A}} : \mathbf{A} \to \mathbf{A}$ *is surjective.*

PROOF: For any $y \in \mathbf{A}$ we have $\mathrm{id}_\mathbf{A}(y) = y$. $\square$

Proposition Schema 3.1.13 (Z)**.** *For any classes* $\mathbf{A}$, $\mathbf{B}$, $\mathbf{C}$, $\mathbf{F}$ *and* $\mathbf{G}$, *the following is a theorem:*

   If $\mathbf{F} : \mathbf{A} \twoheadrightarrow \mathbf{B}$ *and* $\mathbf{G} : \mathbf{B} \twoheadrightarrow \mathbf{C}$, *then* $\mathbf{G} \circ \mathbf{F} : \mathbf{A} \twoheadrightarrow \mathbf{C}$.

PROOF:
$\langle 1 \rangle 1$. LET: $c \in \mathbf{C}$
$\langle 1 \rangle 2$. PICK $b \in \mathbf{B}$ such that $\mathbf{G}(b) = c$.
$\langle 1 \rangle 3$. PICK $a \in \mathbf{A}$ such that $\mathbf{F}(a) = b$.
$\langle 1 \rangle 4$. $(\mathbf{G} \circ \mathbf{F})(a) = c$
$\square$

Proposition 3.1.14 (Z)**.** *Let* $F : A \to B$. *There exists* $H : B \to A$ *(a* right inverse*) such that* $F \circ H = \mathrm{id}_B$ *if and only if* F *maps* A *onto* B.

PROOF:
$\langle 1 \rangle 1$. If F has a right inverse then F is surjective.
   $\langle 2 \rangle 1$. ASSUME: F has a right inverse $H : B \to A$.
   $\langle 2 \rangle 2$. LET: $y \in B$
   $\langle 2 \rangle 3$. $F(H(y)) = y$
   $\langle 2 \rangle 4$. There exists $x \in A$ such that $F(x) = y$
$\langle 1 \rangle 2$. If F is surjective then F has a right inverse.
   $\langle 2 \rangle 1$. ASSUME: F is surjective.
   $\langle 2 \rangle 2$. PICK a function $H : B \to A$ such that $H \subseteq F^{-1}$
      PROOF: Proposition 3.1.6.
   $\langle 2 \rangle 3$. $F \circ H = \mathrm{id}_B$
      $\langle 3 \rangle 1$. LET: $y \in B$
      $\langle 3 \rangle 2$. $(y, H(y)) \in F^{-1}$
      $\langle 3 \rangle 3$. $F(H(y)) = y$
$\square$

3.1.3 Bijections

Definition 3.1.15 (Bijection)**.** Let $\mathbf{F} : \mathbf{A} \to \mathbf{B}$. Then $\mathbf{F}$ is *bijective* or a *bijection*, $\mathbf{F} : \mathbf{A} \approx \mathbf{B}$, iff it is injective and surjective.

Proposition Schema 3.1.16 (Z)**.** *For any class* $\mathbf{A}$, *the following is a theorem:*
   The identity function $\mathrm{id}_\mathbf{A} : \mathbf{A} \approx \mathbf{A}$ *is a bijection.*

PROOF: Proposition 3.1.8 and 3.1.12. $\square$

Proposition Schema 3.1.17 (Z)**.** *For any classes* $\mathbf{A}$, $\mathbf{B}$ *and* $\mathbf{F}$, *the following is a theorem:*
   If $\mathbf{F} : \mathbf{A} \approx \mathbf{B}$ *then* $\mathbf{F}^{-1} : \mathbf{B} \approx \mathbf{A}$.

PROOF:
$\langle 1 \rangle 1$. $\mathbf{F}^{-1} : \mathbf{B} \to \mathbf{A}$
   $\langle 2 \rangle 1$. LET: $b \in \mathbf{B}$

⟨2⟩2. PICK $a \in \mathbf{A}$ such that $\mathbf{F}(a) = b$.
  PROOF: Since $\mathbf{F}$ is surjective.
⟨2⟩3. $(b, a) \in \mathbf{F}^{-1}$
⟨2⟩4. If $(b, a') \in \mathbf{F}^{-1}$ then $a' = a$.
  ⟨3⟩1. LET: $a' \in \mathbf{A}$ such that $(b, a') \in \mathbf{F}^{-1}$
  ⟨3⟩2. $\mathbf{F}(a') = \mathbf{F}(a)$
  ⟨3⟩3. $a' = a$
    PROOF: Since $\mathbf{F}$ is injective.
⟨1⟩2. $\mathbf{F}^{-1}$ is injective.
  ⟨2⟩1. LET: $x, y \in \mathbf{B}$
  ⟨2⟩2. ASSUME: $\mathbf{F}^{-1}(x) = \mathbf{F}^{-1}(y)$
  ⟨2⟩3. $x = y$
    PROOF: $x = \mathbf{F}(\mathbf{F}^{-1}(x)) = \mathbf{F}(\mathbf{F}^{-1}(y)) = y$.
⟨1⟩3. $\mathbf{F}^{-1}$ is surjective.
  PROOF: For all $a \in \mathbf{A}$ we have $\mathbf{F}^{-1}(\mathbf{F}(a)) = a$.
□

Proposition Schema 3.1.18 (Z). *For any classes* $\mathbf{A}$, $\mathbf{B}$, $\mathbf{C}$, $\mathbf{F}$ *and* $\mathbf{G}$, *the following is a theorem:*
  If $\mathbf{F} : \mathbf{A} \approx \mathbf{B}$ *and* $\mathbf{G} : \mathbf{B} \approx \mathbf{C}$ *then* $\mathbf{G} \circ \mathbf{F} : \mathbf{A} \approx \mathbf{C}$.

PROOF: Propositions 3.1.9 and 3.1.13. □

3.1.4 Restrictions

Definition 3.1.19 (Restriction). Let $\mathbf{F} : \mathbf{A} \to \mathbf{B}$. Let $\mathbf{C} \subseteq \mathbf{A}$. The *restriction* of $\mathbf{F}$ to $\mathbf{C}$, denoted $\mathbf{F} \upharpoonright \mathbf{C}$, is the function

$$\mathbf{F} \upharpoonright \mathbf{C} : \mathbf{C} \to \mathbf{B}$$
$$(\mathbf{F} \upharpoonright \mathbf{C})(x) = \mathbf{F}(x) \qquad\qquad (x \in \mathbf{C})$$

3.1.5 Images

Definition 3.1.20 (Image). Let $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ and $\mathbf{C} \subseteq \mathbf{A}$. The *image* of $\mathbf{C}$ under $\mathbf{F}$ is the class
$$\mathbf{F}(\mathbf{C}) := \{\mathbf{F}(x) \mid x \in \mathbf{C}\} \ .$$

Proposition Schema 3.1.21 (Z). *For any classes* $\mathbf{F}$, $\mathbf{A}$ *and* $\mathbf{B}$, *the following is a theorem.*
  If $\mathbf{F} : \mathbf{A} \to \mathbf{B}$, *then for any subset* $S \subseteq \mathbf{A}$, *the class* $\mathbf{F}(S)$ *is a set.*

PROOF: By an Axiom of Replacement. □

Proposition Schema 3.1.22 (Z). *For any classes* $\mathbf{A}$, $\mathbf{B}$, $\mathbf{C}$ *and* $\mathbf{F}$, *the following is a theorem:*
  Assume $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ *and* $\mathbf{C} \subseteq \mathcal{P}\mathbf{A}$. *Then*

$$\mathbf{F}\left(\bigcup \mathbf{C}\right) = \{y \mid \exists X \in \mathbf{C}.y \in \mathbf{F}(X)\}$$

PROOF:
$$y \in \mathbf{F}\left(\bigcup \mathbf{C}\right) \Leftrightarrow \exists x \in \bigcup \mathbf{C}.y = \mathbf{F}(x)$$
$$\Leftrightarrow \exists x.\exists X.X \in \mathbf{C} \wedge x \in X \wedge y = \mathbf{F}(x)$$
$$\Leftrightarrow \exists X \in \mathbf{C}.y \in \mathbf{F}(X) \qquad\qquad \square$$

Proposition Schema 3.1.23 (Z). *For any classes* $\mathbf{A}$, $\mathbf{B}$, $\mathbf{C}$, $\mathbf{D}$ *and* $\mathbf{F}$, *the following is a theorem:*

   Assume $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ *and* $\mathbf{C}, \mathbf{D} \subseteq \mathbf{A}$. *Then*

$$\mathbf{F}(\mathbf{C} \cup \mathbf{D}) = \mathbf{F}(\mathbf{C}) \cup \mathbf{F}(\mathbf{D}) \ .$$

PROOF:
$$y \in \mathbf{F}(\mathbf{C} \cup \mathbf{D}) \Leftrightarrow \exists x \in \mathbf{C} \cup \mathbf{D}.y = \mathbf{F}(x)$$
$$\Leftrightarrow \exists x \in \mathbf{C}.y = \mathbf{F}(x) \vee \exists x \in \mathbf{D}.y = \mathbf{F}(x)$$
$$\Leftrightarrow y \in \mathbf{F}(\mathbf{C}) \cup \mathbf{F}(\mathbf{D}) \qquad\qquad \square$$

Proposition 3.1.24 (Z). *For any classes* $\mathbf{F}$, $\mathbf{A}$, $\mathbf{B}$, $\mathbf{C}$ *and* $\mathbf{D}$, *the following is a theorem:*

   Assume $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ *and* $\mathbf{C}, \mathbf{D} \subseteq \mathbf{A}$. *Then*

$$\mathbf{F}(\mathbf{A} \cap \mathbf{B}) \subseteq \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B}) \ .$$

   Equality holds if $\mathbf{F}$ *is injective.*

PROOF:
$\langle 1 \rangle 1.$ $\mathbf{F}(\mathbf{A} \cap \mathbf{B}) \subseteq \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B})$
   $\langle 2 \rangle 1.$ LET: $y \in \mathbf{F}(\mathbf{A} \cap \mathbf{B})$
   $\langle 2 \rangle 2.$ PICK $x \in \mathbf{A} \cap \mathbf{B}$ such that $y = \mathbf{F}(x)$
   $\langle 2 \rangle 3.$ $y \in \mathbf{F}(\mathbf{A})$
      PROOF: Since $x \in \mathbf{A}$.
   $\langle 2 \rangle 4.$ $y \in \mathbf{F}(\mathbf{B})$
      PROOF: Since $x \in \mathbf{B}$.
$\langle 1 \rangle 2.$ If $\mathbf{F}$ is injective then $\mathbf{F}(\mathbf{A} \cap \mathbf{B}) = \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B})$.
   $\langle 2 \rangle 1.$ ASSUME: $\mathbf{F}$ is injective.
   $\langle 2 \rangle 2.$ LET: $y \in \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B})$
   $\langle 2 \rangle 3.$ PICK $x \in \mathbf{A}$ such that $y = \mathbf{F}(x)$
   $\langle 2 \rangle 4.$ PICK $x' \in \mathbf{B}$ such that $y = \mathbf{F}(x')$
   $\langle 2 \rangle 5.$ $x = x'$
      PROOF: $\langle 2 \rangle 1$
   $\langle 2 \rangle 6.$ $x \in \mathbf{A} \cap \mathbf{B}$
   $\langle 2 \rangle 7.$ $y \in \mathbf{F}(\mathbf{A} \cap \mathbf{B})$
$\square$

Proposition Schema 3.1.25 (Z). *For any classes* $\mathbf{F}$, $\mathbf{A}$, $\mathbf{B}$, *and* $\mathbf{C}$, *the following is a theorem:*

   Let $\mathbf{F} : \mathbf{A} \rightarrow \mathbf{B}$ *and* $\mathbf{C} \subseteq \mathcal{P}\mathbf{A}$. *Then*

$$\mathbf{F}\left(\bigcap \mathbf{C}\right) \subseteq \bigcap\{\mathbf{F}(X) \mid X \in \mathbf{A}\} \ .$$

Equality holds if $\mathbf{F}$ *is injective and* $\mathbf{A}$ *is nonempty.*

PROOF:

$\langle 1 \rangle 1.$ $\mathbf{F}\left(\bigcap \mathbf{A}\right) \subseteq \bigcap\{\mathbf{F}(X) \mid X \in \mathbf{A}\}$

   $\langle 2 \rangle 1.$ LET: $y \in \mathbf{F}(\bigcap \mathbf{A})$

   $\langle 2 \rangle 2.$ PICK $x \in \bigcap \mathbf{A}$ such that $y = \mathbf{F}(x)$

   $\langle 2 \rangle 3.$ LET: $X \in \mathbf{A}$

   $\langle 2 \rangle 4.$ $x \in X$

   $\langle 2 \rangle 5.$ $y \in \mathbf{F}(X)$

$\langle 1 \rangle 2.$ If $\mathbf{F}$ is injective then $\mathbf{F}\left(\bigcap \mathbf{A}\right) = \bigcap\{\mathbf{F}(X) \mid X \in \mathbf{A}\}$

   $\langle 2 \rangle 1.$ ASSUME: $\mathbf{F}$ is injective.

   $\langle 2 \rangle 2.$ ASSUME: $\mathbf{A}$ is nonempty.

   $\langle 2 \rangle 3.$ LET: $y \in \bigcap\{\mathbf{F}(X) \mid X \in \mathbf{A}\}$

   $\langle 2 \rangle 4.$ PICK $X_0 \in \mathbf{A}$

   $\langle 2 \rangle 5.$ PICK $x \in X_0$ such that $(x, y) \in \mathbf{F}$

   $\langle 2 \rangle 6.$ $x \in \bigcap \mathbf{A}$

      $\langle 3 \rangle 1.$ LET: $X \in \mathbf{A}$

      $\langle 3 \rangle 2.$ PICK $x' \in X$ such that $(x', y) \in \mathbf{F}$.

      $\langle 3 \rangle 3.$ $x = x'$

         PROOF: $\langle 2 \rangle 1$

      $\langle 3 \rangle 4.$ $x \in X$

   $\langle 2 \rangle 7.$ $y \in \mathbf{F}(\bigcap \mathbf{A})$

$\square$

Proposition 3.1.26 (Z). *For any classes* $\mathbf{A}$, $\mathbf{B}$, $\mathbf{C}$, $\mathbf{D}$ *and* $\mathbf{F}$, *the following is a theorem:*

   Assume $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ *and* $\mathbf{C}, \mathbf{D} \subseteq \mathbf{A}$. *Then*

$$\mathbf{F}(\mathbf{C}) - \mathbf{F}(\mathbf{D}) \subseteq \mathbf{F}(\mathbf{C} - \mathbf{D}) \ .$$

Equality holds if $\mathbf{F}$ *is injective.*

PROOF:

$\langle 1 \rangle 1.$ $\mathbf{F}(\mathbf{C}) - \mathbf{F}(\mathbf{D}) \subseteq \mathbf{F}(\mathbf{A} - \mathbf{B})$

   $\langle 2 \rangle 1.$ LET: $y \in \mathbf{F}(\mathbf{A}) - \mathbf{F}(\mathbf{B})$

   $\langle 2 \rangle 2.$ PICK $x \in \mathbf{A}$ such that $y = \mathbf{F}(x)$

   $\langle 2 \rangle 3.$ $x \notin \mathbf{B}$

   $\langle 2 \rangle 4.$ $x \in \mathbf{A} - \mathbf{B}$

   $\langle 2 \rangle 5.$ $y \in \mathbf{F}(\mathbf{A} - \mathbf{B})$

$\langle 1 \rangle 2.$ If $\mathbf{F}$ is injective then $\mathbf{F}(\mathbf{A}) - \mathbf{F}(\mathbf{B}) = \mathbf{F}(\mathbf{A} - \mathbf{B})$

   $\langle 2 \rangle 1.$ ASSUME: $\mathbf{F}$ is injective.

   $\langle 2 \rangle 2.$ LET: $y \in \mathbf{F}(\mathbf{A} - \mathbf{B})$

   $\langle 2 \rangle 3.$ PICK $x \in \mathbf{A} - \mathbf{B}$ such that $y = \mathbf{F}(x)$

   $\langle 2 \rangle 4.$ $y \in \mathbf{F}(\mathbf{A})$

   $\langle 2 \rangle 5.$ $y \notin \mathbf{F}(\mathbf{B})$

⟨3⟩1. ASSUME: for a contradiction $y \in \mathbf{F}(\mathbf{B})$
⟨3⟩2. PICK $x' \in \mathbf{B}$ such that $y = \mathbf{F}(x')$
⟨3⟩3. $x = x'$
   PROOF: ⟨2⟩1
⟨3⟩4. $x \in \mathbf{B}$
⟨3⟩5. Q.E.D.
   PROOF: This contradicts ⟨2⟩3.

□

3.1.6   Inverse Images

Definition 3.1.27 (Inverse Image). Let $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ and $\mathbf{C} \subseteq \mathbf{B}$. Then the *inverse image* of $\mathbf{C}$ under $\mathbf{F}$ is

$$\mathbf{F}^{-1}(\mathbf{C}) = \{x \in \mathbf{A} \mid \mathbf{F}(x) \in \mathbf{C}\} \ .$$

Proposition Schema 3.1.28 (Z). *For any classes* $\mathbf{A}$, $\mathbf{B}$, $\mathbf{C}$ *and* $\mathbf{F}$, *the following is a theorem:*
   Assume $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ *and* $\mathbf{C} \subseteq \mathcal{P}\mathbf{B}$. *Then*

$$\mathbf{F}^{-1}\left(\bigcap \mathbf{C}\right) = \bigcap \{\mathbf{F}^{-1}(X) \mid X \in \mathbf{C}\} \ .$$

PROOF:
$$\begin{aligned}
x \in \mathbf{F}^{-1}\left(\bigcap \mathbf{C}\right) &\Leftrightarrow \mathbf{F}(x) \in \bigcap \mathbf{C} \\
&\Leftrightarrow \forall X \in \mathbf{C}.\mathbf{F}(x) \in X \\
&\Leftrightarrow \forall X \in \mathbf{C}.x \in \mathbf{F}^{-1}(X) \qquad \square
\end{aligned}$$

Proposition Schema 3.1.29 (Z). *For any classes* $\mathbf{A}$, $\mathbf{B}$, $\mathbf{C}$, $\mathbf{D}$ *and* $\mathbf{F}$, *the following is a theorem:*
   Assume $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ *and* $\mathbf{C}, \mathbf{D} \subseteq \mathbf{B}$. *Then*

$$\mathbf{F}^{-1}(\mathbf{C} - \mathbf{D}) = \mathbf{F}^{-1}(\mathbf{C}) - \mathbf{F}^{-1}(\mathbf{D}) \ .$$

PROOF:
$$\begin{aligned}
x \in \mathbf{F}^{-1}(\mathbf{C} - \mathbf{D}) &\Leftrightarrow \mathbf{F}(x) \in \mathbf{C} - \mathbf{D} \\
&\Leftrightarrow \mathbf{F}(x) \in \mathbf{C} \wedge \mathbf{F}(x) \notin \mathbf{D} \\
&\Leftrightarrow x \in \mathbf{F}^{-1}(\mathbf{C}) \wedge x \in \mathbf{F}^{-1}(\mathbf{D}) \\
&\Leftrightarrow x \in \mathbf{F}^{-1}(\mathbf{C}) - \mathbf{F}^{-1}(\mathbf{D}) \qquad \square
\end{aligned}$$

3.1.7   Function Sets

Proposition 3.1.30 (ZFC). *For any classes* $\mathbf{B}$ *and* $\mathbf{F}$, *the following is a theorem:*
   Let A *be a set. If* $\mathbf{F} : A \to \mathbf{B}$ *then* $\mathbf{F}$ *is a set.*

PROOF: By an Axiom of Replacement, we have $R = \{\mathbf{F}(x) \mid x \in A\}$ is a set. Hence $\mathbf{F}$ is a set since $\mathbf{F} \subseteq A \times R$. □

Definition 3.1.31 (Dependent Product Class)**.** Let I be a set and let $\mathbf{H}(i)$ be a class for all $i \in I$. We write $\prod_{i \in I} \mathbf{H}(i)$ for the class of all functions $f : I \to \bigcup_{i \in I} \mathbf{H}(i)$ such that $\forall i \in I.f(i) \in \mathbf{H}(i)$.

We write $\mathbf{B}^I$ for $\prod_{i \in I} \mathbf{B}$ where $\mathbf{B}$ does not depend on I.

Proposition Schema 3.1.32 (ZFC)**.** *Let I be a set. Let $H(i)$ be a set for every $i \in I$. Then $\prod_{i \in I} \mathbf{H}(i)$ is a set.*

PROOF:

$\langle 1 \rangle 1$. $\{\mathbf{H}(i) \mid i \in I\}$ is a set.
  PROOF: By an Axiom of Replacement.
$\langle 1 \rangle 2$. $\bigcup_{i \in I} \mathbf{H}(i)$ is a set.
$\langle 1 \rangle 3$. $\prod_{i \in I} \mathbf{H}(i)$ is a set.
  PROOF: It is a subset of $\mathcal{P}\left(I \times \bigcup_{i \in I} \mathbf{H}(i)\right)$.
□

Proposition 3.1.33 (Z)**.** *Let I be a set. Let $H(i)$ be a set for all $i \in I$. If $\forall i \in I.H(i) \neq \emptyset$ then $\prod_{i \in I} H(i) \neq \emptyset$.*

PROOF:

$\langle 1 \rangle 1$. ASSUME: $\forall i \in I.H(i) \neq \emptyset$
$\langle 1 \rangle 2$. LET: $R = \{(i, x) \mid i \in I, x \in H(i)\}$
$\langle 1 \rangle 3$. PICK a function $f : I \to \bigcup_{i \in I} H(i)$ such that $f \subseteq R$
  PROOF: Proposition 3.1.6.
$\langle 1 \rangle 4$. $f \in \prod_{i \in I} H(i)$
□

3.2 Equinumerosity

Definition 3.2.1 (Equinumerous)**.** Sets A and B are *equinumerous*, $A \approx B$, iff there exists a bijection between A and B.

3.3 Domination

Definition 3.3.1 (Dominate)**.** A set A is *dominated* by a set B, $A \preccurlyeq B$, iff there exists an injection $A \to B$.

Proposition 3.3.2 (Z)**.** *Given sets A and B, if $A \neq \emptyset$ or $B = \emptyset$, then we have $A \preccurlyeq B$ iff there exists a surjective function $B \to A$.*

PROOF:

$\langle 1 \rangle 1$. If $A \preccurlyeq B$ and $A \neq \emptyset$ then there exists a surjective function $B \to A$.
  $\langle 2 \rangle 1$. ASSUME: $f : A \to B$ be injective.
  $\langle 2 \rangle 2$. PICK $a \in A$
  $\langle 2 \rangle 3$. LET: $g : B \to A$ be the function defined by $g(b) = f^{-1}(b)$ if $b \in \operatorname{ran} f$, and $g(b) = a$ otherwise.

$\langle 2 \rangle 4$. g is surjective.

$\langle 1 \rangle 2$. If there exists a surjective function $B \to A$ then $A \preccurlyeq B$.

$\langle 2 \rangle 1$. ASSUME: there exists a surjective function $g : B \to A$

$\langle 2 \rangle 2$. $\forall a \in A. \exists b \in B. g(b) = a$

$\langle 2 \rangle 3$. Choose a function $f : A \to B$ such that $\forall a \in A. g(f(a)) = a$

$\langle 2 \rangle 4$. f is injective.

☐

Chapter 4

Equivalence Relations

Definition 4.0.1 (Equivalence Relation)**.** An *equivalence relation* on a class $\mathbf{A}$ is a binary relation on $\mathbf{A}$ that is reflexive, symmetric and transitive.

Proposition 4.0.2 (Z)**.** *Equinumerosity is an equivalence relation on the class of all sets.*

PROOF: Propositions 3.1.16, 3.1.17, 3.1.18. □

Definition 4.0.3 (Respects)**.** Let $\mathbf{R}$ be an equivalence relation on $\mathbf{A}$ and $\mathbf{F} : \mathbf{A} \to \mathbf{B}$. Then $\mathbf{F}$ *respects* $\mathbf{A}$ iff, whenever $(x, y) \in \mathbf{R}$, then $\mathbf{F}(x) = \mathbf{F}(y)$.

Definition 4.0.4 (Equivalence Class)**.** Let $\mathbf{R}$ be an equivalence relation on $\mathbf{A}$ and $a \in \mathbf{A}$. The *equivalence class* of *a modulo* $\mathbf{R}$ is

$$[a]_{\mathbf{R}} := \{x \mid a\mathbf{R}x\} \ .$$

Proposition Schema 4.0.5 (Z)**.** *For any classes $\mathbf{A}$ and $\mathbf{R}$, the following is a theorem.*

Assume $\mathbf{R}$ be an equivalence relation on $\mathbf{A}$. Let $a, b \in \mathbf{A}$. Then $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$ if and only if $a\mathbf{R}b$.

PROOF:
⟨1⟩1. If $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$ then $a\mathbf{R}b$.
   ⟨2⟩1. ASSUME: $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$
   ⟨2⟩2. $b\mathbf{R}b$
     PROOF: Reflexivity
   ⟨2⟩3. $b \in [b]_{\mathbf{R}}$
   ⟨2⟩4. $b \in [a]_{\mathbf{R}}$
   ⟨2⟩5. $a\mathbf{R}b$
⟨1⟩2. If $a\mathbf{R}b$ then $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$.
   ⟨2⟩1. For all $x, y \in \mathbf{A}$, if $x\mathbf{R}y$ then $[y]_{\mathbf{R}} \subseteq [x]_{\mathbf{R}}$
     ⟨3⟩1. LET: $x, y \in \mathbf{A}$
     ⟨3⟩2. ASSUME: $x\mathbf{R}y$

⟨3⟩3. LET: $t \in [y]_{\mathbf{R}}$
⟨3⟩4. $y\mathbf{R}t$
⟨3⟩5. $x\mathbf{R}t$
  PROOF: Transitivity, ⟨3⟩2, ⟨3⟩4.
⟨3⟩6. $t \in [x]_{\mathbf{R}}$
⟨2⟩2. ASSUME: $a\mathbf{R}b$
⟨2⟩3. $[b]_{\mathbf{R}} \subseteq [a]_{\mathbf{R}}$
  PROOF: ⟨2⟩1, ⟨2⟩2.
⟨2⟩4. $b\mathbf{R}a$
  PROOF: Symmetry, ⟨2⟩2.
⟨2⟩5. $[a]_{\mathbf{R}} \subseteq [b]_{\mathbf{R}}$
  PROOF: ⟨2⟩1, ⟨2⟩4.
⟨2⟩6. $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$
  PROOF: ⟨2⟩3, ⟨2⟩5.
☐

Definition 4.0.6 (Partition). A *partition* Π of a set A is a set of nonempty subsets of A that is disjoint and exhaustive, i.e.

1. no two different sets in Π have any common elements, and

2. each element of A is in some set in Π.

Definition 4.0.7. Let R be an equivalence relation on a set A. The *quotient set A/R* is the set of all equivalence classes.

Theorem 4.0.8 (Z). *Let A be a set and* **B** *a class. Let R be an equivalence relation on A and $F : A \to \mathbf{B}$. Then F respects R if and only if there exists $\hat{F} : A/R \to \mathbf{B}$ such that*

$$\forall a \in A.\hat{F}([a]_R) = F(a) \ .$$

In this case, $\hat{F}$ is unique.

PROOF:
⟨1⟩1. If F respects R then there exists $\hat{F} : A/R \to \mathbf{B}$ such that $\forall a \in A.\hat{F}([a]_R) = F(a)$.
  ⟨2⟩1. ASSUME: F respects R.
  ⟨2⟩2. LET: $\hat{F} = \{([a]_R, F(a)) \mid a \in A\}$
  ⟨2⟩3. $\hat{F}$ is a function.
    ⟨3⟩1. ASSUME: $a, a' \in A$ and $[a]_R = [a']_R$
        PROVE: $F(a) = F(a')$
    ⟨3⟩2. $(a, a') \in R$
      PROOF: Proposition 4.0.5.
    ⟨3⟩3. $F(a) = F(a')$
      PROOF: ⟨2⟩1
  ⟨2⟩4. dom $\hat{F} = A/R$
  ⟨2⟩5. ran $\hat{F} \subseteq \mathbf{B}$

$\langle 2 \rangle$6. $\forall a \in A.\hat{F}([a]_R) = F(a)$

$\langle 1 \rangle$2. If there exists $\hat{F} : A/R \to \mathbf{B}$ such that $\forall a \in A.\hat{F}([a]_R) = F(a)$ then F respects R.

$\quad \langle 2 \rangle$1. ASSUME: $\hat{F} : A/R \to \mathbf{B}$ and $\forall a \in A.\hat{F}([a]_R) = F(a)$

$\quad \langle 2 \rangle$2. LET: $a, a' \in A$

$\quad \langle 2 \rangle$3. ASSUME: $(a, a') \in R$

$\quad \langle 2 \rangle$4. $[a]_R = [a']_R$

$\quad\quad$ PROOF: Proposition 4.0.5.

$\quad \langle 2 \rangle$5. $F(a) = F(a')$

$\quad\quad$ PROOF: $\langle 2 \rangle$1

$\langle 1 \rangle$3. If $G, H : A/R \to \mathbf{B}$ and $\forall a \in A.G([a]_R) = H([a]_R)$ then $G = H$.

$\square$

Proposition 4.0.9 (Z). *Let R be an equivalence relation on a set A. Then A/R is a partition of A.*

PROOF:

$\langle 1 \rangle$1. Every member of A/R is nonempty.

$\quad$ PROOF: Since $a \in [a]_R$ by reflexivity.

$\langle 1 \rangle$2. No two different sets in A/R have any common elements.

$\quad \langle 2 \rangle$1. LET: $[a]_R, [b]_R \in A/R$

$\quad \langle 2 \rangle$2. LET: $c \in [a]_R \cap [b]_R$

$\quad\quad$ PROVE: $[a]_R = [b]_R$

$\quad \langle 2 \rangle$3. aRc

$\quad\quad$ PROOF: $\langle 2 \rangle$2

$\quad \langle 2 \rangle$4. bRc

$\quad\quad$ PROOF: $\langle 2 \rangle$2

$\quad \langle 2 \rangle$5. cRb

$\quad\quad$ PROOF: Symmetry, $\langle 2 \rangle$4

$\quad \langle 2 \rangle$6. aRb

$\quad\quad$ PROOF: Transitivity, $\langle 2 \rangle$3, $\langle 2 \rangle$5

$\quad \langle 2 \rangle$7. $[a]_R = [b]_R$

$\quad\quad$ PROOF: Proposition 4.0.5, $\langle 2 \rangle$6

$\langle 1 \rangle$3. Each element of A is in some set in A/R.

$\quad$ PROOF: Since $a \in [a]_R$ by reflexivity.

$\square$

Proposition 4.0.10 (Z). *For any partition P of a set A, there exists a unique equivalence relation R on A such that $A/R = P$, namely xRy iff $\exists X \in P(x \in X \land y \in X)$.*

PROOF: Easy. $\square$

Definition 4.0.11 (Natural Map). Let A be a set and R an equivalence relation on A. The *natural map $A \to A/R$* is the function that maps $a \in A$ to $[a]_R$.

Chapter 5

Ordering Relations

5.1   Partial Orders

Definition 5.1.1 (Partial Ordering). Let **A** be a class. A *partial ordering* on **A** is a relation **R** on **A** that is reflexive, antisymmetric and transitive.

We often write $\leq$ for a partial ordering, and then write $x < y$ for $x \leq y \wedge x \neq y$.

Proposition Schema 5.1.2 (Z). *For any classes* **A** *and* **R**, *the following is a theorem:*

If **R** *is a partial order on* **A** *then so is* $\mathbf{R}^{-1}$.

PROOF:
$\langle 1 \rangle 1.$ $\mathbf{R}^{-1}$ is reflexive.
    PROOF: Proposition 2.2.8.
$\langle 1 \rangle 2.$ $\mathbf{R}^{-1}$ is antisymmetric.
    PROOF: Proposition 2.2.12.
$\langle 1 \rangle 3.$ $\mathbf{R}^{-1}$ is transitive.
    $\langle 2 \rangle 1.$ ASSUME: $x\mathbf{R}^{-1}y$ and $y\mathbf{R}^{-1}z$
    $\langle 2 \rangle 2.$ $y\mathbf{R}x$ and $z\mathbf{R}y$
    $\langle 2 \rangle 3.$ $z\mathbf{R}x$
        PROOF: Since **R** is transitive.
    $\langle 2 \rangle 4.$ $x\mathbf{R}^{-1}z$
☐

Proposition Schema 5.1.3 (Z). *For any classes* **A**, **B**, **F** *and* **R**, *the following is a theorem:*

Assume **R** *is a partial order on* **B** *and* $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ *is injective. Define* **S** *on* **A** *by* $x\mathbf{S}y$ *iff* $\mathbf{F}(x)\mathbf{R}\mathbf{F}(y)$. *Then* **S** *is a partial order on* **A**.

PROOF:
$\langle 1 \rangle 1.$ **S** is reflexive.
    PROOF: For any $x \in \mathbf{A}$ we have $\mathbf{F}(x)\mathbf{R}\mathbf{F}(x)$.

$\langle 1 \rangle 2$. **S** is antisymmetric.
   $\langle 2 \rangle 1$. LET: $x, y \in \mathbf{A}$
   $\langle 2 \rangle 2$. ASSUME: $x\mathbf{S}y$ and $y\mathbf{S}x$
   $\langle 2 \rangle 3$. $\mathbf{F}(x)\mathbf{R}\mathbf{F}(y)$ and $\mathbf{F}(y)\mathbf{R}\mathbf{F}(x)$
   $\langle 2 \rangle 4$. $\mathbf{F}(x) = \mathbf{F}(y)$
     PROOF: **R** is antisymmetric.
   $\langle 2 \rangle 5$. $x = y$
$\langle 1 \rangle 3$. **S** is transitive.
$\square$

Corollary Schema 5.1.3.1 (Z). *For any classes* **A**, **B** *and* **R**, *the following is a theorem:*

Assume **R** *be a partial order on* **A** *and* $\mathbf{B} \subseteq \mathbf{A}$. *Then* $\mathbf{R} \cap \mathbf{B}^2$ *is a partial order on* **B**.

Definition 5.1.4 (Partially Ordered Set). A *partially ordered set* or *poset* is a pair $(A, \leq)$ where A is a set and $\leq$ is a partial ordering on A. We often write just A for $(A, \leq)$.

If $(A, \leq)$ is a poset and $B \subseteq A$ we write just B for the poset $(B, \leq \cap B^2)$.

Definition 5.1.5 (Strictly Monotone). Let $(A, <_A)$ and $(B, <_B)$ be posets. A function $f : A \to B$ is *strictly monotone* iff, whenever $x <_A y$, then $f(x) <_B f(y)$.

Definition 5.1.6 (Least). Let $\leq$ be a partial order on **A**. An element $m \in \mathbf{A}$ is *least* iff for all $x \in \mathbf{A}$ we have $m \leq x$.

Proposition 5.1.7 (Z). *A partial order has at most one least element.*

PROOF: If m and m' are least then $m \leq m'$ and $m' \leq m$, so $m = m'$. $\square$

Definition 5.1.8 (Greatst). Let $\leq$ be a partial order on **A**. An element $m \in \mathbf{A}$ is *greatest* iff for all $x \in A$ we have $x \leq m$.

Proposition 5.1.9 (Z). *A poset has at most one greatest element.*

PROOF: If m and m' are greatest then $m \leq m'$ and $m' \leq m$, so $m = m'$. $\square$

Definition 5.1.10 (Upper Bound). Let $\leq$ be a partial ordering on **A** and $\mathbf{B} \subseteq \mathbf{A}$. Let $u \in \mathbf{A}$. Then u is an *upper bound* for **B** iff $\forall x \in \mathbf{B}.x \leq u$.

Definition 5.1.11 (Lower Bound). Let $\leq$ be a partial ordering on **A** and $\mathbf{B} \subseteq \mathbf{A}$. Let $l \in \mathbf{A}$. Then l is a *lower bound* for **B** iff $\forall x \in \mathbf{B}.l \leq x$.

Definition 5.1.12 (Bounded Above). Let $\leq$ be a partial ordering on **A** and $\mathbf{B} \subseteq \mathbf{A}$. Then **B** is *bounded above* iff it has an upper bound.

Definition 5.1.13 (Bounded Below). Let $\leq$ be a partial ordering on **A** and $\mathbf{B} \subseteq \mathbf{A}$. Then **B** is *bounded below* iff it has a lower bound.

Definition 5.1.14 (Least Upper Bound)**.** Let $\leq$ be a partial ordering on **A** and **B** $\subseteq$ **A**. Let $s \in$ **A**. Then s is the *least upper bound* or *supremum* of **B** iff s is an upper bound for **B** and, for every upper bound u for **B**, we have $s \leq u$.

Definition 5.1.15 (Greatest Lower Bound)**.** Let $\leq$ be a partial ordering on **A** and **B** $\subseteq$ **A**. Let $i \in$ **A**. Then i is the *greatest lower bound* or *infimum* of **B** iff i is a lower bound for **B** and, for every lower bound l for **B**, we have $i \leq l$.

Definition 5.1.16 (Complete)**.** A partial order is *complete* iff every nonempty subset bounded above has a supremum, and every nonempty subset bounded below has an infimum.

Definition 5.1.17 (Order Isomorphism)**.** Let A and B be posets. An *order isomorphism* between A and B, $f : A \cong B$, is a bijection $f : A \approx B$ such that f and f^{-1} are monotone.

Theorem 5.1.18 (Knaster Fixed-Point Theorem (Z))**.** *Let A be a complete poset with a greatest and least element. Let $\phi : A \to A$ be monotone. Then there exists $a \in A$ such that $\phi(a) = a$.*

PROOF:
$\langle 1 \rangle 1$. LET: $B = \{x \in A \mid x \leq \phi(x)\}$
$\langle 1 \rangle 2$. LET: $a = \sup B$
  PROOF: B is nonempty because the least element of A is in B, and it is bounded above by the greatest element of A.
$\langle 1 \rangle 3$. For all $b \in B$ we have $b \leq \phi(a)$
  $\langle 2 \rangle 1$. LET: $b \in B$
  $\langle 2 \rangle 2$. $b \leq \phi(b)$
  $\langle 2 \rangle 3$. $b \leq a$
  $\langle 2 \rangle 4$. $\phi(b) \leq \phi(a)$
  $\langle 2 \rangle 5$. $b \leq \phi(a)$
$\langle 1 \rangle 4$. $a \leq \phi(a)$
$\langle 1 \rangle 5$. $\phi(a) \leq \phi(\phi(a))$
$\langle 1 \rangle 6$. $\phi(a) \in B$
$\langle 1 \rangle 7$. $\phi(a) \leq a$
$\langle 1 \rangle 8$. $\phi(a) = a$
□

Definition 5.1.19 (Dense)**.** Let $\leq$ be a partial order on **A** and **B** $\subseteq$ **A**. Then **B** is *dense* iff, for all $x, y \in$ **A**, if $x < y$ then there exists $z \in$ **B** such that $x < z < y$.

Proposition 5.1.20 (Z)**.** *Let A be a complete poset with no least element. Let $B \subseteq A$ be dense. Let $\theta : A \to A$ be a monotone map that is the identity on B. Then $\theta = \mathrm{id}_A$.*

PROOF:
$\langle 1 \rangle 1$. LET: $a \in A$
    PROVE: $\theta(a) = a$

$\langle 1 \rangle 2$. LET: $S(a) = \{b \in B \mid b < a\}$
$\langle 1 \rangle 3$. $S(a)$ is nonempty and bounded above.
   $\langle 2 \rangle 1$. $S(a)$ is nonempty.
      $\langle 3 \rangle 1$. PICK $a_1 < a$
        PROOF: Since a is not least.
      $\langle 3 \rangle 2$. There exists $b \in B$ such that $a_1 < b < a$.
   $\langle 2 \rangle 2$. $S(a)$ is bounded above by a.
$\langle 1 \rangle 4$. $\sup S(a) \leq a$
$\langle 1 \rangle 5$. $\sup S(a) = a$
   $\langle 2 \rangle 1$. ASSUME: for a contradiction $\sup S(a) < a$
   $\langle 2 \rangle 2$. PICK $b \in B$ such that $\sup S(a) < b < a$
   $\langle 2 \rangle 3$. $b \in S(a)$
   $\langle 2 \rangle 4$. Q.E.D.
   PROOF: This contradicts the fact that $\sup S(a) < b$.
$\langle 1 \rangle 6$. For all $b \in S(a)$ we have $b \leq \theta(a)$
   $\langle 2 \rangle 1$. LET: $b \in S(a)$
   $\langle 2 \rangle 2$. $b < a$
   $\langle 2 \rangle 3$. $\theta(b) \leq \theta(a)$
   $\langle 2 \rangle 4$. $b \leq \theta(a)$
   PROOF: $\theta(b) = b$
$\langle 1 \rangle 7$. $a \leq \theta(a)$
  PROOF: Since $a = \sup S(a)$ and $\theta(a)$ is an upper bound for $S(a)$.
$\langle 1 \rangle 8$. $a \not< \theta(a)$
   $\langle 2 \rangle 1$. ASSUME: for a contradiction $a < \theta(a)$.
   $\langle 2 \rangle 2$. PICK $b \in B$ such that $a < b < \theta(a)$
   $\langle 2 \rangle 3$. $\theta(a) \leq \theta(b) = b$
   $\langle 2 \rangle 4$. Q.E.D.
   PROOF: This contradicts the fact that $b < \theta(a)$.
$\langle 1 \rangle 9$. $\theta(a) = a$
☐

Theorem 5.1.21 (Z). *Let A and P be complete posets with no least or greatest element. Let B be dense in A and Q be dense in P. Every order isomorphism $\phi : B \cong Q$ extends uniquely to an order isomorphism $A \cong P$.*

PROOF:
$\langle 1 \rangle 1$. For $a \in A$, let $S(a) = \{b \in B \mid b < a\}$.
$\langle 1 \rangle 2$. Define $\overline{\phi} : A \to P$ by $\overline{\phi}(a) = \sup \phi(S(a))$.
   $\langle 2 \rangle 1$. $\phi(S(a))$ is nonempty.
      $\langle 3 \rangle 1$. PICK $a_1 < a$
        PROOF: Since a is not least.
      $\langle 3 \rangle 2$. PICK $b \in B$ such that $a_1 < b < a$.
      $\langle 3 \rangle 3$. $\phi(b) \in \phi(S(a))$
   $\langle 2 \rangle 2$. $\phi(S(a))$ is bounded above.
      $\langle 3 \rangle 1$. PICK $a_2 > a$
        PROOF: Since a is not greatest.
      $\langle 3 \rangle 2$. PICK $b \in B$ such that $a < b < a_2$

⟨3⟩3. $\phi(b)$ is an upper bound for $\phi(S(a))$.

⟨1⟩3. $\overline{\phi}$ is monotone.

PROOF: If $a \leq a'$ then $S(a) \subseteq S(a')$ and so $\overline{\phi}(a) \leq \overline{\phi}(a')$.

⟨1⟩4. $\overline{\phi}$ extends ϕ.

  ⟨2⟩1. LET: $b \in B$

    PROVE: $\phi(b) = \sup \phi(S(b))$

  ⟨2⟩2. $\phi(b)$ is an upper bound for $\phi(S(b))$

  ⟨2⟩3. LET: u be any upper bound for $\phi(S(b))$

    PROVE: $\phi(b) \leq u$

  ⟨2⟩4. ASSUME: for a contradiction $u < \phi(b)$

  ⟨2⟩5. PICK $q \in Q$ such that $u < q < \phi(b)$

  ⟨2⟩6. PICK $b' \in B$ such that $\phi(b') = q$

  ⟨2⟩7. $b' < b$

  ⟨2⟩8. $b' \in S(b)$

  ⟨2⟩9. $q = \phi(b') \leq u$

  ⟨2⟩10. Q.E.D.

    PROOF: This is a contradiction.

⟨1⟩5. LET: $\psi = \phi^{-1}$

⟨1⟩6. LET: $\overline{\psi} : P \rightarrow A$ be the function $\overline{\psi}(p) = \sup\{\psi(q) \mid q \in Q, q < p\}$

⟨1⟩7. $\overline{\psi}$ is monotone and extends ψ

  PROOF: Similar.

⟨1⟩8. $\overline{\psi} \circ \overline{\phi} : A \rightarrow A$ is monotone and the identity on B.

⟨1⟩9. $\overline{\psi} \circ \overline{\phi} = \mathrm{id}_A$

  PROOF: Proposition 5.1.20.

⟨1⟩10. $\overline{\phi} \circ \overline{\psi} = \mathrm{id}_B$

  PROOF: Proposition 5.1.20.

⟨1⟩11. If $\phi^* : A \cong P$ is any order isomorphism that extends ϕ then $\phi^* = \overline{\phi}$.

  ⟨2⟩1. LET: $a \in A$

    PROVE: $\phi^*(a) = \sup \phi(S(a))$

  ⟨2⟩2. $\phi^*(a)$ is an upper bound for $\phi(S(a))$

  ⟨2⟩3. LET: u be any upper bound for $\phi(S(a))$

    PROVE: $\phi^*(a) \leq u$

  ⟨2⟩4. ASSUME: for a contradiction $u < \phi^*(a)$

  ⟨2⟩5. PICK $q \in Q$ such that $u < q < \phi^*(a)$

  ⟨2⟩6. PICK $b \in B$ such that $q = \phi(b)$

  ⟨2⟩7. $b < a$

  ⟨2⟩8. $b \in S(a)$

  ⟨2⟩9. $q = \phi(b) \leq u$

  ⟨2⟩10. Q.E.D.

    PROOF: This is a contradiction.

☐

Definition 5.1.22 (Initial Segment)**.** Let $\leq$ be a partial order on **A** and $t \in A$. The *initial segment* up to t is the class

$$\operatorname{seg} t := \{x \in \mathbf{A} \mid x < t\} \ .$$

Definition 5.1.23 (Lexicographic Ordering)**.** Let $\mathbf{R}$ be a partial order on $\mathbf{A}$ and $\mathbf{S}$ a partial order on $\mathbf{B}$. The *lexicographic ordering* $\leq$ on $\mathbf{A} \times \mathbf{B}$ is defined by:

$$(a, b) \leq (a', b') \Leftrightarrow (a\mathbf{R}a' \wedge a \neq a') \vee (a = a' \wedge b\mathbf{S}b') \ .$$

Proposition Schema 5.1.24 (Z)**.** *For any classes* $\mathbf{A}$, $\mathbf{B}$, $\mathbf{R}$ *and* $\mathbf{S}$, *the following is a theorem:*

   If $\mathbf{R}$ *is a partial order on* $\mathbf{A}$ *and* $\mathbf{S}$ *is a partial order on* $\mathbf{B}$ *then the lexicographic ordering on* $\mathbf{A} \times \mathbf{B}$ *is a partial order.*

PROOF:

⟨1⟩1. LET: $\leq$ be the lexicographic ordering on $\mathbf{A} \times \mathbf{B}$

⟨1⟩2. $\leq$ is reflexive.

   PROOF: For any $a \in \mathbf{A}$ and $b \in \mathbf{B}$ we have $a = a$ and $b\mathbf{S}b$, so $(a, b) \leq (a, b)$.

⟨1⟩3. $\leq$ is antisymmetric.

   ⟨2⟩1. ASSUME: $(a, b) \leq (a', b')$ and $(a', b') \leq (a, b)$

   ⟨2⟩2. $(a\mathbf{R}a' \wedge a \neq a') \vee (a = a' \wedge b\mathbf{S}b')$

   ⟨2⟩3. $(a'\mathbf{R}a \wedge a' \neq a) \vee (a' = a \wedge b\mathbf{S}b')$

   ⟨2⟩4. CASE: $a = a'$

      PROOF: Then $b\mathbf{S}b'$ and $b'\mathbf{S}b$ hence $b = b'$ and $(a, b) = (a', b')$.

   ⟨2⟩5. CASE: $a \neq a'$

      PROOF: Then $a\mathbf{R}a'$ and $a'\mathbf{R}a$ hence $a = a'$ which is a contradiction.

⟨1⟩4. $\leq$ is transitive.

   ⟨2⟩1. ASSUME: $(a_1, b_1) \leq (a_2, b_2) \leq (a_3, b_3)$

   ⟨2⟩2. $(a_1\mathbf{R}a_2 \wedge a_1 \neq a_2) \vee (a_1 = a_2 \wedge b_1\mathbf{S}b_2)$

   ⟨2⟩3. $(a_2\mathbf{R}a_3 \wedge a_2 \neq a_3) \vee (a_2 = a_3 \wedge b_2\mathbf{S}b_3)$

   ⟨2⟩4. CASE: $a_1\mathbf{R}a_2, a_1 \neq a_2, a_2\mathbf{R}a_3, a_2 \neq a_3$

      ⟨3⟩1. $a_1\mathbf{R}a_3$

         PROOF: Since $\mathbf{R}$ is transitive.

      ⟨3⟩2. $a_1 \neq a_3$

         PROOF: If $a_1 = a_3$ then $a_1\mathbf{R}a_2$ and $a_2\mathbf{R}a_1$ so $a_1 = a_2$ which is a contradiction.

   ⟨2⟩5. CASE: $a_1\mathbf{R}a_2, a_1 \neq a_2, a_2 = a_3, b_2\mathbf{S}b_3$

      PROOF: Then $a_1\mathbf{R}a_3$ and $a_1 \neq a_3$.

   ⟨2⟩6. CASE: $a_1 = a_2, b_1\mathbf{S}b_2, a_2\mathbf{R}a_3, a_2 \neq a_3$

      PROOF: Then $a_1\mathbf{R}a_3$ and $a_1 \neq a_3$.

   ⟨2⟩7. CASE: $a_1 = a_2, b_1\mathbf{S}b_2, a_2 = a_3, b_2\mathbf{S}b_3$

      PROOF: Then $a_1 = a_3$ and $b_1\mathbf{S}b_3$.

□

5.2   Linear Orders

Definition 5.2.1 (Linear Ordering)**.** Let $\mathbf{A}$ be a class. A *linear ordering* or *total ordering* on $\mathbf{A}$ is a partial ordering $\leq$ on $\mathbf{A}$ that is *total*, i.e.

$$\forall x, y \in \mathbf{A}.x \leq y \vee y \leq x$$

We often use the symbol $<$ for a linear ordering, and then write $x < y$ for $(x, y) \in <$.

Proposition Schema 5.2.2 (Trichotomy (Z)). *For any classes* **A** *and* $\leq$*, the following is a theorem:*

Assume $\leq$ *be a linear ordering on* **A**. *For any* $x, y \in$ **A***, exactly one of $x < y$, $x = y$, $y < x$ holds.*

PROOF: Immediate from definitions. □

Proposition Schema 5.2.3 (Z). *For any classes* **A** *and* $<$*, the following is a theorem:*

Let $<$ *be a transitive relation on* **A** *that satisfies trichotomy. Define* $\leq$ *on* **A** *by $x \leq y$ iff $x < y$ or $x = y$. Then $\leq$ is a linear ordering on* **A** *and $x < y$ iff $x \leq y$ and $x \neq y$.*

PROOF:
$\langle 1 \rangle 1.$ $\leq$ is reflexive.
  PROOF: By definition we have $\forall x \in$ **A**$.x \leq x$.
$\langle 1 \rangle 2.$ $\leq$ is antisymmetric.
  $\langle 2 \rangle 1.$ ASSUME: $x \leq y$ and $y \leq x$
  $\langle 2 \rangle 2.$ $x < y$ or $x = y$
  $\langle 2 \rangle 3.$ $y < x$ or $y = x$
  $\langle 2 \rangle 4.$ We cannot have $x < y$ and $y < x$
    PROOF: Trichotomy.
  $\langle 2 \rangle 5.$ $x = y$
$\langle 1 \rangle 3.$ $\leq$ is transitive.
  $\langle 2 \rangle 1.$ ASSUME: $x \leq y$ and $y \leq z$
  $\langle 2 \rangle 2.$ $x < y$ or $x = y$
  $\langle 2 \rangle 3.$ $y < z$ or $y = z$
  $\langle 2 \rangle 4.$ CASE: $x < y$ and $y < z$
    PROOF: Then $x < z$ by transitivity, so $x \leq z$.
  $\langle 2 \rangle 5.$ CASE: $x = y$
    PROOF: Then we have $y \leq z$ and so $x \leq z$.
  $\langle 2 \rangle 6.$ CASE: $y = z$
    PROOF: Then we have $x \leq y$ and so $x \leq z$.
$\langle 1 \rangle 4.$ $\leq$ is total.
  PROOF: Immediate from trichotomy.
□

Proposition Schema 5.2.4 (Z). *For any classes* **A** *and* **R***, the following is a theorem:*

If **R** *is a linear ordering on* **A** *then* $\mathbf{R}^{-1}$ *is also a linear ordering on* **A**.

PROOF:
$\langle 1 \rangle 1.$ $\mathbf{R}^{-1}$ is a partial order on **A**.
  PROOF: Proposition 5.1.2.
$\langle 1 \rangle 2.$ $\mathbf{R}^{-1}$ is total.

$\langle 2 \rangle 1$. Let: $x, y \in \mathbf{A}$
$\langle 2 \rangle 2$. $x\mathbf{R}y$ or $y\mathbf{R}x$.
$\langle 2 \rangle 3$. $y\mathbf{R}^{-1}x$ or $x\mathbf{R}^{-1}y$.
$\square$

Proposition Schema 5.2.5 (Z)**.** *For any classes* $\mathbf{A}$, $\mathbf{B}$, $\mathbf{F}$, $\mathbf{R}$, $\mathbf{S}$, *the following is a theorem:*

Assume $\mathbf{R}$ *is a linear order on* $\mathbf{A}$, $\mathbf{S}$ *is a partial order on* $\mathbf{B}$, *and* $\mathbf{F} : \mathbf{A} \to \mathbf{B}$. *If* $\mathbf{F}$ *is strictly monotone then it is injective.*

Proof:
$\langle 1 \rangle 1$. Let: $x, y \in \mathbf{A}$
$\langle 1 \rangle 2$. Assume: $x \neq y$
     Prove: $\mathbf{F}(x) \neq \mathbf{F}(y)$
$\langle 1 \rangle 3$. Assume: w.l.o.g. $x\mathbf{R}y$
   Proof: $\mathbf{R}$ is total.
$\langle 1 \rangle 4$. $\mathbf{F}(x)\mathbf{SF}(y)$ and $\mathbf{F}(x) \neq \mathbf{F}(y)$
   Proof: $\mathbf{F}$ is strictly monotone.
$\square$

Proposition Schema 5.2.6 (Z)**.** *For any classes* $\mathbf{A}$, $\mathbf{B}$, $\leq$, $\preccurlyeq$ *and* $\mathbf{F}$, *the following is a theorem:*

Assume $\leq$ *is a linear order on* $\mathbf{A}$ *and* $\preccurlyeq$ *is a linear order on* $\mathbf{B}$. *Assume* $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ *and* $\mathbf{F}$ *is strictly monotone. For all* $x, y \in \mathbf{A}$, *if* $\mathbf{F}(x) \prec \mathbf{F}(y)$ *then* $x < y$.

Proof:
$\langle 1 \rangle 1$. $\mathbf{F}(x) \neq \mathbf{F}(y)$ and $\mathbf{F}(y) \not\prec \mathbf{F}(x)$
   Proof: Trichotomy.
$\langle 1 \rangle 2$. $x \neq y$ and $y \not< x$
   Proof: $\mathbf{F}$ is strictly monotone.
$\langle 1 \rangle 3$. $x < y$
   Proof: Trichotomy.
$\square$

Corollary Schema 5.2.6.1 (Z)**.** *For any classes* $\mathbf{A}$, $\mathbf{B}$, $\leq$, $\preccurlyeq$ *and* $\mathbf{F}$, *the following is a theorem:*

Assume $\leq$ *is a linear order on* $\mathbf{A}$ *and* $\preccurlyeq$ *is a linear order on* $\mathbf{B}$. *Assume* $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ *and* $\mathbf{F}$ *is strictly monotone. Then* $\mathbf{F}$ *is an order isomorphism.*

Proposition Schema 5.2.7 (Z)**.** *For any classes* $\mathbf{A}$, $\mathbf{B}$, $\mathbf{F}$ *and* $\mathbf{S}$, *the following is a theorem:*

Assume $\mathbf{S}$ *is a linear order on* $\mathbf{B}$ *and* $\mathbf{F} : \mathbf{A} \rightarrowtail \mathbf{B}$. *Define* $\mathbf{R}$ *on* $\mathbf{A}$ *by* $x\mathbf{R}y$ *if and only if* $\mathbf{F}(x)\mathbf{SF}(y)$. *Then* $\mathbf{R}$ *is a linear order on* $\mathbf{A}$.

Proof:
$\langle 1 \rangle 1$. $\mathbf{R}$ is a partial order on $\mathbf{A}$.
   Proof: Proposition 5.1.3.

⟨1⟩2. **R** is total.
   PROOF: For all $x, y \in \mathbf{A}$ we have $\mathbf{F}(x)\mathbf{SF}(y)$ or $\mathbf{F}(y)\mathbf{SF}(x)$.
□

Corollary Schema 5.2.7.1 (Z). *For any classes* **A**, **B** *and* **R**, *the following is a theorem:*

   Assume **R** *be a linear order on* **A** *and* $\mathbf{B} \subseteq \mathbf{A}$. *Then* $\mathbf{R} \cap \mathbf{B}^2$ *is a linear order on* **B**.

Proposition Schema 5.2.8 (Z). *For any classes* **A**, **B**, **R** *and* **S**, *the following is a theorem:*

   Assume **R** *is a linear order on* **A** *and* **S** *is a linear order on* **B**. *Then the lexicographic ordering is a linear order on* $\mathbf{A} \times \mathbf{B}$.

PROOF:
⟨1⟩1. LET: $\leq$ be the lexicographic order on $\mathbf{A} \times \mathbf{B}$
⟨1⟩2. $\leq$ is a partial order.
   PROOF: Proposition 5.1.24.
⟨1⟩3. $\leq$ is total.
   ⟨2⟩1. LET: $a, a' \in \mathbf{A}$ and $b, b' \in \mathbf{B}$
   ⟨2⟩2. CASE: $a\mathbf{R}a'$ and $a \neq a'$
     PROOF: Then $(a, b) \leq (a', b')$.
   ⟨2⟩3. CASE: $a = a'$
     PROOF: We have $b\mathbf{S}b'$ or $b'\mathbf{S}b$, so $(a, b) \leq (a', b')$ or $(a', b') \leq (a, b)$.
   ⟨2⟩4. CASE: $a'\mathbf{R}a$ and $a \neq a'$
     PROOF: Then $(a', b') \leq (a, b)$.
□

5.3   Well Orderings

Definition 5.3.1 (Well Ordering). A *well ordering* on a class **A** is a well-founded linear ordering on **A**.

Proposition 5.3.2 (Z). *Let S be a well ordering of the set B and $f : A \to B$ a function. Define R on A by xRy if and only if $F(x)SF(y)$. Then R well orders A.*

PROOF:
⟨1⟩1. R linearly orders A.
   PROOF: Proposition 5.2.7.
⟨1⟩2. Every nonempty subset of A has a least element.
   ⟨2⟩1. LET: C be a nonempty subset of A.
   ⟨2⟩2. LET: y be the least element of $f(C)$.
   ⟨2⟩3. PICK $x \in C$ such that $f(x) = y$.
   ⟨2⟩4. x is least in C.
□

Proposition Schema 5.3.3 (Z)**.** *For any classes* **A***,* **B** *and* **R***, the following is a theorem:*

   Assume **R** *well orders* **B** *and* **A** $\subseteq$ **B***. Then* **R** $\cap$ **A**^{2} *well orders* **A***.*

Proof:
$\langle 1 \rangle$1. Let: **R**$'$ = **R** $\cap$ **A**^{2}
$\langle 1 \rangle$2. **R**$'$ linearly orders **A**.
   Proof: Corollary 5.2.7.1.
$\langle 1 \rangle$3. **R**$'$ is well founded.
   Proof: Proposition 2.4.4.
$\Box$

Proposition Schema 5.3.4 (ZFC)**.** *For any classes* **A***,* **B***,* **F** *and* **S***, the following is a theorem:*

   Assume **S** *well orders* **B** *and* **F** : **A** $\rightarrowtail$ **B***. Define* **R** *on* **A** *by* x**R**y *if and only if* **F**(x)**SF**(y)*. Then* **R** *well orders* **A***.*

Proof:
$\langle 1 \rangle$1. **R** linearly orders **A**.
   Proof: Proposition 5.2.7.
$\langle 1 \rangle$2. For all $t \in$ **A** we have $\{x \in \mathbf{A} \mid x\mathbf{R}t \wedge x \neq t\}$ is a set.
   $\langle 2 \rangle$1. Let: $t \in$ **A**
   $\langle 2 \rangle$2. Let: $S = \{y \in \mathbf{B} \mid y\mathbf{SF}(t) \wedge y \neq \mathbf{F}(t)\}$
   $\langle 2 \rangle$3. Let: $P(x, y)$ be the property $\mathbf{F}(y) = x$
   $\langle 2 \rangle$4. For all $x \in S$ there exists at most one y such that $P(x, y)$
      Proof: **F** is injective.
   $\langle 2 \rangle$5. Let: $T = \{y \mid \exists x \in S.P(x, y)\}$
      Proof: Axiom of Replacement.
   $\langle 2 \rangle$6. $T = \{x \in \mathbf{A} \mid x\mathbf{R}t \wedge x \neq t\}$
$\langle 1 \rangle$3. Every nonempty subset of **A** has a least element.
   $\langle 2 \rangle$1. Let: S be a nonempty subset of **A**.
   $\langle 2 \rangle$2. $\mathbf{F}(S)$ is a nonempty subset of **B**
      Proof: Axiom of Replacement.
   $\langle 2 \rangle$3. Let: y be the least element of $\mathbf{F}(S)$.
   $\langle 2 \rangle$4. Pick $x \in S$ such that $\mathbf{F}(x) = y$.
   $\langle 2 \rangle$5. x is least in S.
$\Box$

Proposition 5.3.5 (Z)**.** *For any well ordered sets A and B, the lexicographic order well orders $A \times B$.*

Proof:
$\langle 1 \rangle$1. $A \times B$ is linearly ordered.
   Proof: Proposition 5.2.8.
$\langle 1 \rangle$2. Every nonempty subset of $A \times B$ has a least element.
   $\langle 2 \rangle$1. Let: S be a nonempty subset of $A \times B$.
   $\langle 2 \rangle$2. Let: a be the least element of $\{x \in A \mid \exists y \in B.(x, y) \in S\}$.
   $\langle 2 \rangle$3. Let: b be the least element of $\{y \in B \mid (a, y) \in S\}$.

$\langle 2 \rangle 4.$ (a, b) is least in S.

□

Definition 5.3.6 (End Extension)**.** Let A and B be well ordered sets. Then B is an *end extension* of A iff $A \subseteq B$ and:

- Whenever $x, y \in A$ then $x \leq_A y$ iff $x \leq_B y$.

- Whenever $x \in A$ and $y \in B - A$ then $x < y$.

Theorem 5.3.7 (Z)**.** *Let $\leq$ be a linear ordering on A. Assume that, for any $B \subseteq A$ such that $\forall t \in A. \operatorname{seg} t \subseteq B \Rightarrow t \in B$, we have $B = A$. Then $\leq$ is a well ordering on A.*

PROOF:
$\langle 1 \rangle 1.$ LET: $C \subseteq A$ be nonempty.
$\langle 1 \rangle 2.$ LET: $B = \{t \in A \mid \forall x \in C. t < x\}$
$\langle 1 \rangle 3.$ $B \cap C = \emptyset$
$\langle 1 \rangle 4.$ $B \neq A$
$\langle 1 \rangle 5.$ PICK $t \in A$ such that $\operatorname{seg} t \subseteq B$ and $t \notin B$
$\langle 1 \rangle 6.$ t is least in C.

□

Proposition Schema 5.3.8 (Z)**.** *For any classes $\mathbf{A}$, $\mathbf{B}$, $\mathbf{F}$, $\mathbf{G}$, $\leq$ and $\preccurlyeq$, the following is a theorem:*
Assume $\leq$ well orders $\mathbf{A}$ and $\preccurlyeq$ well orders $\mathbf{B}$. Assume $\mathbf{F}$ and $\mathbf{G}$ are order isomorphisms between $\mathbf{A}$ and $\mathbf{B}$. Then $\mathbf{F} = \mathbf{G}$.

PROOF:
$\langle 1 \rangle 1.$ For all $x \in \mathbf{A}$, if $\forall t < x. \mathbf{F}(t) = \mathbf{G}(t)$, then $\mathbf{F}(x) = \mathbf{G}(x)$
　$\langle 2 \rangle 1.$ LET: $x \in \mathbf{A}$
　$\langle 2 \rangle 2.$ ASSUME: $\forall t < x. \mathbf{F}(t) = \mathbf{G}(t)$
　$\langle 2 \rangle 3.$ $\mathbf{F}(\operatorname{seg} x) = \mathbf{G}(\operatorname{seg} x)$
　$\langle 2 \rangle 4.$ $\mathbf{F}(x)$ is the least element of $\mathbf{B} - \mathbf{F}(\operatorname{seg} x)$
　$\langle 2 \rangle 5.$ $\mathbf{G}(x)$ is the least element of $\mathbf{B} - \mathbf{G}(\operatorname{seg} x)$
　$\langle 2 \rangle 6.$ $\mathbf{F}(x) = \mathbf{G}(x)$
$\langle 1 \rangle 2.$ $\forall x \in \mathbf{A}. \mathbf{F}(x) = \mathbf{G}(x)$
　PROOF: Transfinite induction.

□

Theorem 5.3.9 (ZFC)**.** *Let A and B be well ordered sets. Then one of the following holds: $A \cong B$; there exists $b \in B$ such that $A \cong \operatorname{seg} b$; there exists $a \in A$ such that $\operatorname{seg} a \cong B$.*

PROOF:
$\langle 1 \rangle 1.$ PICK e that is not in A or B.
$\langle 1 \rangle 2.$ LET: $F : A \to B \cup \{e\}$ be the function defined by transfinite recursion thus:
$$F(t) = \begin{cases} \text{the least element of } B - F(\operatorname{seg} t) & \text{if } B - F(\operatorname{seg} t) \neq \emptyset \\ e & \text{if } B - F(\operatorname{seg} t) = \emptyset \end{cases}$$

⟨1⟩3. CASE: $e \in \operatorname{ran} F$
    ⟨2⟩1. LET: t be least such that $F(t) = e$
    ⟨2⟩2. $F \restriction \operatorname{seg} t : \operatorname{seg} t \cong B$
⟨1⟩4. CASE: $\operatorname{ran} F = B$
    PROOF: We have $F : A \cong B$
⟨1⟩5. CASE: $\operatorname{ran} F \subsetneq B$
    ⟨2⟩1. LET: b be the least element of $B - \operatorname{ran} F$
    ⟨2⟩2. $F : A \cong \operatorname{seg} b$
∎

Chapter 6

Ordinal Numbers

6.1 Ordinals

Definition 6.1.1 (Ordinal Number). An *ordinal (number)* is a transitive set α that is *well-ordered by* $\in$; that is, such that $\{(x, y) \in \alpha^2 \mid x \in y \lor x = y\}$ well orders α.

Given $x, y \in \alpha$, we write $x < y$ iff $x \in y$, and $x \leq y$ iff $x \in y$ or $x = y$.

Let **On** be the class of ordinal numbers. For $\alpha, \beta \in$ **On**, we write $\alpha < \beta$ iff $\alpha \in \beta$, and $\alpha \leq \beta$ iff $\alpha < \beta$ or $\alpha = \beta$.

Proposition 6.1.2 (Z). *For any ordinal numbers α and β, if $\alpha \cong \beta$ then $\alpha = \beta$.*

Proof:
$\langle 1 \rangle 1$. Let: $f : \alpha \cong \beta$
$\langle 1 \rangle 2$. For all $x \in \alpha$, if $\forall t < x.f(t) = t$ then $f(x) = x$
  $\langle 2 \rangle 1$. $f(x) \subseteq x$
    $\langle 3 \rangle 1$. Let: $y \in f(x)$
    $\langle 3 \rangle 2$. $y \in \beta$
    $\langle 3 \rangle 3$. Pick $t \in \alpha$ such that $f(t) = y$
      Proof: f is surjective.
    $\langle 3 \rangle 4$. $f(t) \in f(x)$
    $\langle 3 \rangle 5$. $t \in x$
      Proof: Since f is an order isomorphism.
    $\langle 3 \rangle 6$. $f(t) = t$
      Proof: Induction hypothesis.
    $\langle 3 \rangle 7$. $y = t$
    $\langle 3 \rangle 8$. $y \in x$
  $\langle 2 \rangle 2$. $x \subseteq f(x)$
    $\langle 3 \rangle 1$. Let: $t \in x$
    $\langle 3 \rangle 2$. $f(t) \in f(x)$
    $\langle 3 \rangle 3$. $f(t) = t$
    $\langle 3 \rangle 4$. $t \in f(x)$

$\langle 1 \rangle 3.$ $\forall x \in \alpha . f(x) = x$
   Proof: Transfinite induction.
$\langle 1 \rangle 4.$ $\alpha = \beta$
   Proof: Since $\beta = \{f(t) \mid t \in \alpha\} = \{t \mid t \in \alpha\} = \alpha$.
▯

Theorem 6.1.3 (ZFC). *Every well-ordered set is isomorphic to a unique ordinal.*

Proof:
$\langle 1 \rangle 1.$ For any well-ordered set A, there exists an ordinal α such that $A \cong \alpha$.
  $\langle 2 \rangle 1.$ Let: A be a well-ordered set.
  $\langle 2 \rangle 2.$ Define the function E on A by transfinite recursion thus:
$$E(t) = \{E(x) \mid x < t\} \qquad (t \in A) \ .$$
  $\langle 2 \rangle 3.$ Let: $\alpha = \{E(x) \mid x \in A\}$
  $\langle 2 \rangle 4.$ α is an ordinal.
    $\langle 3 \rangle 1.$ α is a transitive set.
      $\langle 4 \rangle 1.$ Let: $x \in y \in \alpha$
      $\langle 4 \rangle 2.$ Pick $t \in A$ such that $y = E(t)$
      $\langle 4 \rangle 3.$ $x \in E(t) = \{E(s) \mid s < t\}$
      $\langle 4 \rangle 4.$ Pick $s < t$ such that $x = E(s)$
      $\langle 4 \rangle 5.$ $x \in \alpha$
    $\langle 3 \rangle 2.$ α is well-ordered by $\in$.
      $\langle 4 \rangle 1.$ Let: $<= \{(x, y) \in \alpha \mid x \in y\}$
      $\langle 4 \rangle 2.$ $<$ is transitive.
        $\langle 5 \rangle 1.$ Let: $x, y, z \in \alpha$ with $x \in y \in z$
        $\langle 5 \rangle 2.$ Pick $t \in A$ such that $z = E(t)$
        $\langle 5 \rangle 3.$ Pick $s \in A$ such that $s < t$ and $y = E(s)$
        $\langle 5 \rangle 4.$ Pick $r \in A$ such that $r < s$ and $x = E(r)$
        $\langle 5 \rangle 5.$ $r < t$
        $\langle 5 \rangle 6.$ $x \in z$
      $\langle 4 \rangle 3.$ $<$ satisfies trichotomy.
        $\langle 5 \rangle 1.$ Let: $x, y \in \alpha$
        $\langle 5 \rangle 2.$ Pick $s, t \in A$ such that $E(s) = x$ and $E(t) = y$
        $\langle 5 \rangle 3.$ Exactly one of $s < t$, $s = t$, $t < s$ holds.
        $\langle 5 \rangle 4.$ Case: $s < t$
          $\langle 6 \rangle 1.$ $x \in y$
          $\langle 6 \rangle 2.$ $x \neq y$ and $y \notin x$
            Proof: Axiom of Regularity.
        $\langle 5 \rangle 5.$ Case: $s = t$
          $\langle 6 \rangle 1.$ $x = y$
          $\langle 6 \rangle 2.$ $x \notin y$ and $y \notin x$
            Proof: Axiom of Regularity.
        $\langle 5 \rangle 6.$ Case: $t < s$
          Proof: Similar to $\langle 5 \rangle 4$.
      $\langle 4 \rangle 4.$ $\leq$ is a linear order on α.
        Proof: Proposition 5.2.3.

$\langle 4 \rangle 5$. Every nonempty subset of α has a least element.

    $\langle 5 \rangle 1$. LET: S be a nonempty subset of α

    $\langle 5 \rangle 2$. LET: $T = \{x \in A \mid E(x) \in S\}$

    $\langle 5 \rangle 3$. LET: t be the least element of T.

        PROVE: $E(t)$ is least in S

    $\langle 5 \rangle 4$. LET: $y \in S$

    $\langle 5 \rangle 5$. PICK $s \in T$ such that $E(s) = y$

    $\langle 5 \rangle 6$. $t \leq s$

    $\langle 5 \rangle 7$. $x \leq y$

$\langle 2 \rangle 5$. E is surjective.

  PROOF: By definition of α.

$\langle 2 \rangle 6$. E is strictly monotone.

  PROOF: If $s < t$ then $E(s) \in E(t)$ by definition of $E(t)$.

$\langle 2 \rangle 7$. Q.E.D.

  PROOF: Corollary 5.2.6.1.

$\langle 1 \rangle 2$. For any ordinals α and β, if $\alpha \cong \beta$ then $\alpha = \beta$.

  PROOF: Proposition 6.1.2.

$\square$

Proposition 6.1.4 (Z). *The class* **On** *is a transitive class. That is, every element of an ordinal is an ordinal.*

PROOF:

$\langle 1 \rangle 1$. LET: α be an ordinal.

$\langle 1 \rangle 2$. LET: $\beta \in \alpha$

$\langle 1 \rangle 3$. β is a transitive set.

  $\langle 2 \rangle 1$. LET: $x \in y \in \beta$

  $\langle 2 \rangle 2$. $y \in \alpha$

    PROOF: α is transitive.

  $\langle 2 \rangle 3$. $x \in \alpha$

    PROOF: α is transitive.

  $\langle 2 \rangle 4$. $x \in \beta$

    PROOF: Since $\{(x, y) \in \alpha^2 \mid x \in y\}$ is transitive.

$\langle 1 \rangle 4$. β is well ordered by $\in$.

  PROOF: By Proposition 5.3.3.

$\square$

Proposition 6.1.5 (ZFC). *Given two ordinal numbers α, β, exactly one of $\alpha \in \beta$, $\alpha = \beta$, $\beta \in \alpha$ holds.*

PROOF:

$\langle 1 \rangle 1$. At most one holds.

  PROOF: Since every ordinal is a transitive set and we never have $\alpha \in \alpha$.

$\langle 1 \rangle 2$. At least one holds.

  $\langle 2 \rangle 1$. Either $\alpha \cong \beta$ or $\exists t \in \beta . \alpha \cong \operatorname{seg} t$ or $\exists t \in \alpha . \operatorname{seg} t \cong \beta$.

  $\langle 2 \rangle 2$. CASE: $\alpha \cong \beta$

    PROOF: Then $\alpha = \beta$ by Proposition 6.1.2.

⟨2⟩3. CASE: There exists $t \in \beta$ such that $\alpha \cong \operatorname{seg} t$

  ⟨3⟩1. t is an ordinal number.
    PROOF: Proposition 6.1.4.

  ⟨3⟩2. $t = \operatorname{seg} t$

    ⟨4⟩1. $t \subseteq \operatorname{seg} t$

      ⟨5⟩1. LET: $s \in t$

      ⟨5⟩2. $s \in \beta$
        PROOF: β is a transitive set.

      ⟨5⟩3. $s \in \operatorname{seg} t$

    ⟨4⟩2. $\operatorname{seg} t \subseteq t$
      PROOF: Immediate from definitions.

  ⟨3⟩3. $\alpha = t$
    PROOF: Proposition 6.1.2.

  ⟨3⟩4. $\alpha \in \beta$

⟨2⟩4. CASE: There exists $t \in \alpha$ such that $\operatorname{seg} t \cong \beta$
  PROOF: $\beta \in \alpha$ similarly.

☐

Proposition 6.1.6 (Z)**.** *Any nonempty set S of ordinal numbers has a least element.*

PROOF:
⟨1⟩1. PICK $\beta \in S$
⟨1⟩2. CASE: $\beta \cap S = \emptyset$
  PROOF: Then β is least in S.
⟨1⟩3. CASE: $\beta \cap S \neq \emptyset$
  PROOF: The least element of $\beta \cap S$ is least in S.
☐

Theorem 6.1.7 (ZFC)**.** *The class **On** is well ordered by $\in$.*

PROOF:
⟨1⟩1. LET: $\mathbf{E} = \{(x, y) \in \mathbf{On}^2 \mid x \in y\}$
⟨1⟩2. $\mathbf{E}$ is transitive.
  PROOF: If $\alpha \in \beta \in \gamma$ then $\alpha \in \gamma$ because every ordinal is a transitive set.
⟨1⟩3. $\mathbf{E}$ satisfies trichotomy.
  PROOF: Proposition 6.1.5.
⟨1⟩4. $\mathbf{E}$ linearly orders **On**.
  PROOF: Proposition 5.2.3.
⟨1⟩5. $\mathbf{E}$ is well founded.
  PROOF: Proposition 2.4.2.
☐

Corollary 6.1.7.1 (Burali-Forti Paradox (ZFC))**.** *The class **On** is a proper class.*

PROOF: If it were a set, it would be a transitive set well-ordered by $\in$, and hence a member of itself, contradicting Proposition 1.5.3.

Proposition 6.1.8 (ZFC). *Any transitive set of ordinal numbers is an ordinal number.*

PROOF: It is well-ordered by $\in$ by Proposition 5.3.3 and Theorem 6.1.7. □

Proposition 6.1.9 (Z). *$\emptyset$ is an ordinal number.*

PROOF: Vacuously, it is a transitive set well-ordered by $\in$.

Definition 6.1.10. We define $0 = \emptyset$.

Proposition 6.1.11 (ZFC). *If A is a set of ordinal numbers then $\bigcup A$ is an ordinal number.*

PROOF:
⟨1⟩1. $\bigcup A$ is a transitive set.
   PROOF: Proposition 1.6.3.
⟨1⟩2. $\bigcup A$ is a set of ordinals.
   PROOF: Proposition 6.1.4.
⟨1⟩3. Q.E.D.
   PROOF: Proposition 6.1.8.
□

Corollary 6.1.11.1 (ZFC). *The poset* **On** *is complete.*

PROOF: For any nonempty set A of ordinals, $\bigcup A$ is its supremum. □

Proposition 6.1.12 (ZFC). *Let α be an ordinal and $S \subseteq \alpha$. Then S is well-ordered by $\in$ and the ordinal of $(S, \in)$ is $\leq \alpha$.*

PROOF:
⟨1⟩1. S is well ordered by $\in$.
⟨1⟩2. LET: β be the ordinal of $(S, \in)$
⟨1⟩3. LET: $E : S \approx \beta$ be the unique isomorphism.
⟨1⟩4. $\forall \gamma \in S . E(\gamma) \leq \gamma$
   ⟨2⟩1. LET: $\gamma \in S$
   ⟨2⟩2. ASSUME: as transfinite induction hypothesis $\forall \delta < \gamma . E(\delta) \leq \delta$
   ⟨2⟩3. $E(\gamma)$ is the least element of β that is greater than $E(\delta)$ for all $\delta < \gamma$
   ⟨2⟩4. γ is greater than $E(\delta)$ for all $\delta < \gamma$
   ⟨2⟩5. $E(\gamma) \leq \gamma$
⟨1⟩5. $\beta \leq \alpha$
   ⟨2⟩1. $\forall \gamma < \beta . \gamma < \alpha$
      ⟨3⟩1. LET: $\gamma < \beta$
      ⟨3⟩2. PICK $\delta \in S$ such that $E(\delta) = \gamma$
      ⟨3⟩3. $\gamma = E(\delta) \leq \delta < \alpha$
□

Proposition 6.1.13 (ZFC). *Let α be a set. Then the following are equivalent.*

   1. α is an ordinal.

2. α is a transitive set and, for all $x, y \in \alpha$, either $x = y$ or $x \in y$ or $y \in x$.

3. α is a transitive set of transitive sets.

PROOF:

⟨1⟩1. $1 \Rightarrow 2$
   PROOF: Immediate from definitions.
⟨1⟩2. $2 \Rightarrow 3$
   ⟨2⟩1. ASSUME: α is a transitive set and, for all $x, y \in \alpha$, either $x = y$ or $x \in y$
                 or $y \in x$
   ⟨2⟩2. LET: $z \in \alpha$
         PROVE: z is transitive.
   ⟨2⟩3. LET: $x \in y \in z$
   ⟨2⟩4. $y \in \alpha$
   ⟨2⟩5. $x \in \alpha$
   ⟨2⟩6. Either $x = z$ or $x \in z$ or $z \in x$
   ⟨2⟩7. $x \neq z$
       PROOF: We cannot have $x \in y \in x$ by the Axiom of Regularity.
   ⟨2⟩8. $z \notin x$
       PROOF: We cannot have $x \in y \in z \in x$ by the Axiom of Regularity.
⟨1⟩3. $3 \Rightarrow 1$
   ⟨2⟩1. LET: x be a transitive set of transitive sets.
   ⟨2⟩2. ASSUME: as $\in$-induction hypothesis that, for all $y \in x$, if y is a transitive
                 set of transitive sets then y is a transitive set of ordinals.
   ⟨2⟩3. Every element of x is an ordinal.
       ⟨3⟩1. LET: $y \in x$
       ⟨3⟩2. y is transitive.
       ⟨3⟩3. Every element of y is transitive.
         PROOF: Since every element of y is an element of x, because x is transi-
         tive.
       ⟨3⟩4. y is an ordinal.
         PROOF: ⟨2⟩2
   ⟨2⟩4. Q.E.D.
       PROOF: Proposition 6.1.8.
☐

Lemma 6.1.14 (Z). *Let A and B be well-ordered sets. If B is an end extension of A then the ordinal of A is $\leq$ the ordinal of B.*

PROOF:

⟨1⟩1. LET: α be the ordinal of A and β the ordinal of B.
⟨1⟩2. LET: $E_A : A \cong \alpha$ and $E_B : B \cong \beta$ be the canonical isomorphisms.
⟨1⟩3. $\forall a \in A . E_A(a) = E_B(a)$
   ⟨2⟩1. LET: $a \in A$
   ⟨2⟩2. ASSUME: as transfinite induction hypothesis $\forall x < a . E_A(x) = E_B(x)$
   ⟨2⟩3. $E_A(a)$ is the least ordinal that is greater than $E_A(x)$ for all $x < a$
   ⟨2⟩4. $E_B(a)$ is the least ordinal that is greater than $E_B(x)$ for all $x < b$

$\langle 2 \rangle 5.$ $\{x \in A \mid x <_A a\} = \{x \in B \mid x <_B a\}$
$\langle 2 \rangle 6.$ $E_A(a) = E_B(a)$
$\langle 1 \rangle 4.$ $\alpha \subseteq \beta$
$\langle 1 \rangle 5.$ $\alpha \leq \beta$

▯

Lemma 6.1.15. *Let $\mathcal{C}$ be a set of well ordered sets such that, for any $A, B \in \mathcal{C}$, we have that one of A and B is an end extension of the other. Let $W = \bigcup \mathcal{C}$ under $x \leq y$ iff there exists $A \in W$ such that $x, y \in A$ and $x \leq y$. Then W is a well ordered set whose ordinal is the supremum of the ordinals of the members of $\mathcal{C}$.*

PROOF:
$\langle 1 \rangle 1.$ $\leq$ is reflexive on W.
  $\langle 2 \rangle 1.$ LET: $x \in W$
  $\langle 2 \rangle 2.$ PICK $A \in W$ such that $x \in A$.
  $\langle 2 \rangle 3.$ $x \leq x$
$\langle 1 \rangle 2.$ $\leq$ is antisymmetric on W.
  $\langle 2 \rangle 1.$ LET: $x, y \in W$
  $\langle 2 \rangle 2.$ ASSUME: $x \leq y$ and $y \leq x$
  $\langle 2 \rangle 3.$ PICK $A \in W$ such that $x, y \in A$ and $x \leq_A y$, and $B \in W$ such that $x, y \in B$ and $y \leq_B x$
  $\langle 2 \rangle 4.$ ASSUME: w.l.o.g. B is an end extension of A
  $\langle 2 \rangle 5.$ $x \leq_B y$ and $y \leq_B x$
  $\langle 2 \rangle 6.$ $x = y$
$\langle 1 \rangle 3.$ $\leq$ is transitive on W.
  $\langle 2 \rangle 1.$ ASSUME: $x \leq y \leq z$
  $\langle 2 \rangle 2.$ PICK $A, B \in W$ such that $x \leq_A y$ and $y \leq_B z$
  $\langle 2 \rangle 3.$ CASE: A is an end extension of B.
    $\langle 3 \rangle 1.$ $x \leq_A y$ and $y \leq_A z$
    $\langle 3 \rangle 2.$ $x \leq_A z$
    $\langle 3 \rangle 3.$ $x \leq z$
  $\langle 2 \rangle 4.$ CASE: B is an end extension of A.
    PROOF: Similar.
$\langle 1 \rangle 4.$ $\leq$ is total on W.
  $\langle 2 \rangle 1.$ LET: $x, y \in W$
  $\langle 2 \rangle 2.$ PICK $A, B \in \mathcal{C}$ such that $x \in A$ and $y \in B$
  $\langle 2 \rangle 3.$ ASSUME: w.l.o.g. B is an end extension of A
  $\langle 2 \rangle 4.$ $x \leq_B y$ or $y \leq_B x$
  $\langle 2 \rangle 5.$ $x \leq_W y$ or $y \leq_W x$
$\langle 1 \rangle 5.$ Every nonempty subset of W has a least element.
  $\langle 2 \rangle 1.$ LET: S be a nonempty subset of W
  $\langle 2 \rangle 2.$ PICK $s \in S$
  $\langle 2 \rangle 3.$ PICK $A \in \mathcal{C}$ such that $s \in A$
  $\langle 2 \rangle 4.$ LET: a be the $\leq_A$-least element of $S \cap A$
       PROVE: a is least in S
  $\langle 2 \rangle 5.$ LET: $x \in S$

  Prove: $a \leq x$

$\langle 2 \rangle 6$. Pick $B \in \mathcal{C}$ such that $x \in B$

$\langle 2 \rangle 7$. Case: A is an end extension of B

  $\langle 3 \rangle 1$. $a \leq_A x$

  $\langle 3 \rangle 2$. $a \leq x$

$\langle 2 \rangle 8$. Case: B is an end extension of A

  $\langle 3 \rangle 1$. Case: $x \in A$

   $\langle 4 \rangle 1$. $a \leq_A x$

   $\langle 4 \rangle 2$. $a \leq x$

  $\langle 3 \rangle 2$. Case: $x \in B - A$

   $\langle 4 \rangle 1$. $a \leq_B x$

   $\langle 4 \rangle 2$. $a \leq x$

$\langle 1 \rangle 6$. For all $A \in \mathcal{C}$, W is an end extension of A.

  $\langle 2 \rangle 1$. For all $x, y \in A$, we have $x \leq_A y$ if and only if $x \leq_W y$

  $\langle 3 \rangle 1$. Let: $x, y \in A$

  $\langle 3 \rangle 2$. If $x \leq_A y$ then $x \leq_W y$

   Proof: Immediate from definitions.

  $\langle 3 \rangle 3$. If $x \leq_W y$ then $x \leq_A y$

   $\langle 4 \rangle 1$. Assume: $x \leq_W y$

   $\langle 4 \rangle 2$. Pick $B \in \mathcal{C}$ such that $x \leq_B y$

   $\langle 4 \rangle 3$. Case: A is an end extension of B

     Proof: Then $x \leq_A y$.

   $\langle 4 \rangle 4$. Case: B is an end extension of A

     Proof: Then $x \leq_A y$.

  $\langle 2 \rangle 2$. For all $x \in A$ and $y \in W - A$ we have $x < y$

  $\langle 3 \rangle 1$. Let: $x \in A$ and $y \in W - A$

  $\langle 3 \rangle 2$. Pick $B \in \mathcal{C}$ such that $y \in B$

  $\langle 3 \rangle 3$. B is an end extension of A

  $\langle 3 \rangle 4$. $x <_B y$

  $\langle 3 \rangle 5$. $x <_W y$

$\langle 1 \rangle 7$. For all $A \in \mathcal{C}$, the ordinal of A is $\leq$ the ordinal of W.

  Proof: Lemma 6.1.14.

$\langle 1 \rangle 8$. For any ordinal α, if for all $A \in \mathcal{C}$ the ordinal of A is $\leq \alpha$, then the ordinal
     of W is $\leq \alpha$.

  $\langle 2 \rangle 1$. Let: α be an ordinal.

  $\langle 2 \rangle 2$. Assume: for all $A \in \mathcal{C}$, the ordinal of A is $\leq \alpha$

  $\langle 2 \rangle 3$. Let: β be the ordinal of W

  $\langle 2 \rangle 4$. Let: $E : W \approx \beta$ be the canonical isomorphism.

  $\langle 2 \rangle 5$. Assume: for a contradiction $\alpha < \beta$

  $\langle 2 \rangle 6$. Let: $a \in W$ be the element with $E(a) = \alpha$

  $\langle 2 \rangle 7$. Pick $A \in \mathcal{C}$ such that $a \in A$

  $\langle 2 \rangle 8$. Let: γ be the ordinal of A and $E_A : A \cong \gamma$ be the canonical isomor-
       phism.

  $\langle 2 \rangle 9$. For all $x \in A$ we have $E_A(x) = E(x)$

   Proof: Transfinite induction on x.

  $\langle 2 \rangle 10$. $E_A(a) = \alpha$

$\langle 2 \rangle 11.$ $\alpha < \gamma$
$\langle 2 \rangle 12.$ Q.E.D.
    PROOF: This contradicts $\langle 2 \rangle 2$.
☐

6.2 Successors

Definition 6.2.1 (Successor)**.** The *successor* of a set a is the set $a^+ := a \cup \{a\}$.

Proposition 6.2.2 (Z)**.** *A set a is a transitive set if and only if*

$$\bigcup (a^+) = a \ .$$

PROOF:
$\langle 1 \rangle 1.$ If a is a transitive set then $\bigcup (a^+) = a$.
    $\langle 2 \rangle 1.$ ASSUME: a is a transitive set.
    $\langle 2 \rangle 2.$ $\bigcup (a^+) \subseteq a$
        $\langle 3 \rangle 1.$ LET: $x \in \bigcup (a^+)$
            PROVE: $x \in a$
        $\langle 3 \rangle 2.$ PICK $y \in a^+$ such that $x \in y$.
        $\langle 3 \rangle 3.$ $y \in a$ or $y = a$.
        $\langle 3 \rangle 4.$ CASE: $y \in a$
          PROOF: Then $x \in a$ because a is a transitive set.
        $\langle 3 \rangle 5.$ CASE: $y = a$
          PROOF: Then $x \in a$ immediately.
    $\langle 2 \rangle 3.$ $a \subseteq \bigcup (a^+)$
      PROOF: Since $a \in a^+$.
$\langle 1 \rangle 2.$ If $\bigcup (a^+) = a$ then a is a transitive set.
    $\langle 2 \rangle 1.$ ASSUME: $\bigcup (a^+) = a$
    $\langle 2 \rangle 2.$ $\bigcup a \subseteq a$
      PROOF:

$$\bigcup a \subseteq \bigcup (a^+) \qquad\qquad \text{(Proposition 1.5.9)}$$
$$= a \qquad\qquad\qquad (\langle 2 \rangle 1)$$

    $\langle 2 \rangle 3.$ a is a transitive set.
      PROOF: Proposition 1.6.2.
☐

Proposition 6.2.3. *For any set a, we have a is a transitive set if and only if a^+ is a transitive set.*

PROOF:
$\langle 1 \rangle 1.$ If a is a transitive set then a^+ is a transitiveset.
    PROOF: If a is a transitive set then $\bigcup (a^+) = a \subseteq a^+$ by Proposition 6.2.2 and
    so a^+ is a transitive set.
$\langle 1 \rangle 2.$ If a^+ is a transitive set then a is a transitive set.
    $\langle 2 \rangle 1.$ ASSUME: a^+ is a transitive set.

$\langle 2 \rangle 2$. LET: $x \in y \in a$
$\langle 2 \rangle 3$. $x \in y \in a^+$
$\langle 2 \rangle 4$. $x \in a^+$
   PROOF: $\langle 2 \rangle 1$
$\langle 2 \rangle 5$. $x \neq a$
   PROOF: From $\langle 2 \rangle 2$ and the Axiom of Regularity.
$\langle 2 \rangle 6$. $x \in a$
□

Definition 6.2.4. We write 0 for $\emptyset$, 1 for $\emptyset^+$, 2 for $\emptyset^{++}$, etc.

Proposition 6.2.5. *For any set A we have $\mathcal{P}A \approx 2^A$.*

PROOF: The function $H : \mathcal{P}A \to 2^A$ defined by $H(S)(a) = \{\emptyset\}$ if $a \in S$ and $\emptyset$ if $a \notin S$ is a bijection. □

Proposition 6.2.6. *For any ordinal number α we have α^+ is an ordinal number.*

PROOF:
$\langle 1 \rangle 1$. α^+ is a transitive set.
  PROOF: Proposition 6.2.3.
$\langle 1 \rangle 2$. α^+ is well-ordered by $\in$.
  $\langle 2 \rangle 1$. For all $x, y, z \in \alpha^+$, if $x \in y \in z$ then $x \in z$
    $\langle 3 \rangle 1$. CASE: $z = \alpha$
     PROOF: Then $x \in \alpha$ since α is a transitive set.
    $\langle 3 \rangle 2$. CASE: $z \in \alpha$
     PROOF: Then $x \in z$ since α is well-ordered by $\in$.
  $\langle 2 \rangle 2$. For all $x, y \in \alpha^+$ we have $x \in y$ or $x = y$ or $y \in x$
    $\langle 3 \rangle 1$. CASE: $x, y \in \alpha$
     PROOF: The result follows because α is well-ordered by $\in$.
    $\langle 3 \rangle 2$. CASE: $x \in \alpha$, $y = \alpha$
     PROOF: Then $x \in y$.
    $\langle 3 \rangle 3$. CASE: $x = \alpha$, $y \in \alpha$
     PROOF: Then $y \in x$.
    $\langle 3 \rangle 4$. CASE: $x = \alpha$, $y = \alpha$
     PROOF: Then $x = y$.
  $\langle 2 \rangle 3$. Every nonempty subset of α^+ has an $\in$-least element.
    $\langle 3 \rangle 1$. LET: $S \subseteq \alpha^+$ be nonempty
    $\langle 3 \rangle 2$. CASE: $S = \{\alpha\}$
     PROOF: α is least in S.
    $\langle 3 \rangle 3$. CASE: $S \neq \{\alpha\}$
      $\langle 4 \rangle 1$. $S - \{\alpha\}$ is a nonempty subset of α
      $\langle 4 \rangle 2$. LET: β be least in $S - \{\alpha\}$
      $\langle 4 \rangle 3$. β is least in S.
□

Proposition 6.2.7. *For ordinals α and β, if $\alpha^+ = \beta^+$ then $\alpha = \beta$.*

PROOF: If $\alpha^+ = \beta^+$ then

$$\alpha = \bigcup(\alpha^+) \qquad\qquad \text{(Proposition 6.2.2)}$$
$$= \bigcup(\beta^+)$$
$$= \beta \qquad\qquad \text{(Proposition 6.2.2)}$$

Proposition 6.2.8. *For ordinals α and β, we have $\alpha < \beta$ if and only if $\alpha^+ < \beta^+$.*

PROOF:

$$\alpha < \beta \Leftrightarrow \alpha^+ \leq \beta$$
$$\Leftrightarrow \alpha^+ < \beta^+ \qquad\qquad □$$

Definition 6.2.9 (Successor Ordinal). An ordinal α is a *successor ordinal* iff $\alpha = \beta^+$ for some β.

Definition 6.2.10 (Limit Ordinal). A *limit ordinal* is an ordinal that is neither 0 nor a successor ordinal.

Proposition 6.2.11. *If λ is a limit ordinal and $\beta < \lambda$ then $\beta^+ < \lambda$.*

PROOF: Since $\beta^+ \leq \lambda$ and $\beta^+ \neq \lambda$. □

6.3   The Well-Ordering Theorem and Zorn's Lemma

Theorem 6.3.1 (Hartogs). *For any set A, there exists an ordinal not dominated by A.*

PROOF:
$\langle 1\rangle 1$. LET: α be the class of all ordinals β such that $\beta \preccurlyeq A$
     PROVE: α is a set.
$\langle 1\rangle 2$. LET: $W = \{(B, R) \mid B \subseteq A, R \text{ is a well ordering on } B\}$
$\langle 1\rangle 3$. α is the class of the ordinals of the elements of W.
  $\langle 2\rangle 1$. For all $(B, R) \in W$, the ordinal of (B, R) is in α.
    $\langle 3\rangle 1$. LET: $(B, R) \in W$
    $\langle 3\rangle 2$. LET: β be the ordinal of (B, R)
    $\langle 3\rangle 3$. LET: $E : B \cong \beta$ be the canonical isomorphism.
    $\langle 3\rangle 4$. LET: $i : B \hookrightarrow A$ be the inclusion
    $\langle 3\rangle 5$. $i \circ E^{-1}$ is an injection $\beta \to A$
    $\langle 3\rangle 6$. $\beta \in \alpha$
  $\langle 2\rangle 2$. For all $\beta \in \alpha$, there exists $(B, R) \in W$ such that β is the ordinal number of (B, R).
    $\langle 3\rangle 1$. LET: $\beta \in \alpha$
    $\langle 3\rangle 2$. PICK an injection $f : \beta \to A$
    $\langle 3\rangle 3$. Define $\leq$ on ran f by $f(x) \leq f(y)$ iff $x \leq y$
    $\langle 3\rangle 4$. $(\text{ran } f, \leq) \in W$
    $\langle 3\rangle 5$. β is the ordinal number of $(\text{ran } f, \leq)$

$\langle 1 \rangle 4.$ α is a set.

    PROOF: By an Axiom of Replacement.

$\langle 1 \rangle 5.$ α is an ordinal.

    PROOF: It is a transitive set of ordinals.

$\langle 1 \rangle 6.$ $\alpha \npreceq A$

    PROOF: Since $\alpha \notin \alpha$.

$\Box$

Theorem 6.3.2 (Numeration Theorem). *Every set is equinumerous with some ordinal.*

PROOF:

$\langle 1 \rangle 1.$ LET: A be any set.

$\langle 1 \rangle 2.$ PICK an ordinal α not dominated by A.

$\langle 1 \rangle 3.$ PICK a choice function G for A.

$\langle 1 \rangle 4.$ PICK $e \notin A$

$\langle 1 \rangle 5.$ LET: $F : \alpha \to A \cup \{e\}$ by transfinite recursion:

$$F(\gamma) = \begin{cases} G(A - F(\{\delta \mid \delta < \gamma\})) & \text{if } A - F(\{\delta \mid \delta < \gamma\}) \neq \emptyset \\ e & \text{if } A - F(\{\delta \mid \delta < \gamma\}) = \emptyset \end{cases}$$

$\langle 1 \rangle 6.$ $e \in \operatorname{ran} F$

    $\langle 2 \rangle 1.$ ASSUME: for a contradiction $e \notin \operatorname{ran} F$

    $\langle 2 \rangle 2.$ F is an injection $\alpha \to A$.

        $\langle 3 \rangle 1.$ LET: $\beta, \gamma \in \alpha$ with $\beta \neq \gamma$

            PROVE: $F(\beta) \neq F(\gamma)$

        $\langle 3 \rangle 2.$ ASSUME: w.l.o.g. $\beta < \gamma$

        $\langle 3 \rangle 3.$ $F(\gamma) \in A - F(\{\delta \mid \delta < \gamma\})$

        $\langle 3 \rangle 4.$ $F(\gamma) \notin F(\{\delta \mid \delta < \gamma\})$

        $\langle 3 \rangle 5.$ $F(\gamma) \neq F(\beta)$

    $\langle 2 \rangle 3.$ Q.E.D.

        PROOF: This contradicts $\langle 1 \rangle 2$.

$\langle 1 \rangle 7.$ LET: δ be least such that $F(\delta) = e$

$\langle 1 \rangle 8.$ $F \restriction \delta : \delta \approx A$

Theorem 6.3.3 (Well-Ordering Theorem). *Any set can be well ordered.*

PROOF:

$\langle 1 \rangle 1.$ PICK an ordinal δ and a bijection $F : A \approx \delta$

$\langle 1 \rangle 2.$ Define $\leq$ on A by $F(x) \leq F(y)$ iff $x \leq y$ for $x, y \in \delta$

$\langle 1 \rangle 3.$ $\leq$ is a well ordering on A.

$\Box$

Theorem 6.3.4 (Zorn's Lemma). *Let $\mathcal{A}$ be a set such that, for every chain $\mathcal{B} \subseteq \mathcal{A}$, we have $\bigcup \mathcal{B} \in \mathcal{A}$. Then $\mathcal{A}$ has a maximal element.*

PROOF:

$\langle 1 \rangle 1.$ PICK a well ordering $<$ on $\mathcal{A}$.

⟨1⟩2. LET: $F : \mathcal{A} \to 2$ be the function defined by transfinite recursion by:
$$F(A) = \begin{cases} 1 & \text{if } A \text{ includes every set } B < A \text{ for which } F(B) = 1 \\ 0 & \text{otherwise} \end{cases}$$
⟨1⟩3. LET: $\mathcal{C} = \{A \in \mathcal{A} \mid F(A) = 1\}$
    PROVE: $\bigcup \mathcal{C}$ is a maximal element of $\mathcal{A}$
⟨1⟩4. For all $A \in \mathcal{A}$, we have $A \in \mathcal{C}$ iff $\forall B < A . B \in \mathcal{C} \Rightarrow B \subseteq A$
⟨1⟩5. $\mathcal{C}$ is a chain.
  ⟨2⟩1. LET: $A, A' \in \mathcal{C}$
  ⟨2⟩2. ASSUME: w.l.o.g. $A \leq A'$
  ⟨2⟩3. $A \subseteq A'$
    PROOF: By ⟨1⟩4
⟨1⟩6. $\bigcup \mathcal{C} \in \mathcal{A}$
⟨1⟩7. $\bigcup \mathcal{C}$ is maximal in $\mathcal{A}$.
  ⟨2⟩1. LET: $A \in \mathcal{A}$ and $\bigcup \mathcal{C} \subseteq A$
  ⟨2⟩2. $A \in \mathcal{C}$
    PROOF: By ⟨1⟩4 since $\forall B \in \mathcal{C} . B \subseteq A$.
  ⟨2⟩3. $A \subseteq \bigcup \mathcal{C}$
  ⟨2⟩4. $A = \bigcup \mathcal{C}$
□

Proposition 6.3.5 (Teichmüller-Tukey Lemma). *Let $\mathcal{A}$ be a nonempty set such that, for every B, we have $B \in \mathcal{A}$ if and only if every finite subset of B is a member of $\mathcal{A}$. Then $\mathcal{A}$ has a maximal element.*

PROOF:
⟨1⟩1. For every chain $\mathcal{B} \subseteq \mathcal{A}$, we have $\bigcup \mathcal{B} \in \mathcal{A}$
  ⟨2⟩1. LET: $\mathcal{B} \subseteq \mathcal{A}$ be a chain.
  ⟨2⟩2. Every finite subset of $\bigcup \mathcal{B}$ is a member of $\mathcal{A}$.
    ⟨3⟩1. LET: C be a finite subset of $\bigcup \mathcal{B}$.
    ⟨3⟩2. PICK $B \in \mathcal{B}$ such that $C \subseteq B$.
    ⟨3⟩3. $B \in \mathcal{A}$
    ⟨3⟩4. Every finite subset of B is in $\mathcal{A}$.
    ⟨3⟩5. $C \in \mathcal{A}$
  ⟨2⟩3. $\bigcup \mathcal{B} \in \mathcal{A}$.
⟨1⟩2. Q.E.D.
  PROOF: Zorn's lemma.
□

Theorem Schema 6.3.6. *For any class $\mathbf{A}$, there exists a class $\mathbf{F}$ such that the following is a theorem:*

    If $\mathbf{A}$ is a proper class of ordinals, then $\mathbf{F} : \mathbf{On} \to \mathbf{A}$ is an order isomorphism.

PROOF:
⟨1⟩1. Define $\mathbf{F} : \mathbf{On} \to \mathbf{A}$ by transfinite recursion as follows: $\mathbf{F}(\alpha)$ is the least element of $\mathbf{A}$ that is different from $\mathbf{F}(\beta)$ for all $\beta < \alpha$.
⟨1⟩2. For all $\alpha, \beta \in \mathbf{On}$, if $\alpha < \beta$ then $\mathbf{F}(\alpha) < \mathbf{F}(\beta)$

PROOF: We have $\mathbf{F}(\alpha) \neq \mathbf{F}(\beta)$ by the definition of $\mathbf{F}(\beta)$, and $\mathbf{F}(\beta) \not< \mathbf{F}(\alpha)$ by the leastness of $\mathbf{F}(\alpha)$.

$\langle 1 \rangle 3$. $\mathbf{F}$ is surjective.

$\quad \langle 2 \rangle 1$. LET: $\alpha \in \mathbf{A}$

$\quad \langle 2 \rangle 2$. ASSUME: as transfinite induction hypothesis $\forall \beta \in \mathbf{A}$, if $\beta < \alpha$ then there exists γ such that $\beta = \mathbf{F}(\gamma)$.

$\quad \langle 2 \rangle 3$. LET: $\gamma = \{\delta \in \mathbf{On} \mid \mathbf{F}(\delta) < \alpha\}$

$\quad \langle 2 \rangle 4$. γ is a set.

$\quad$ PROOF: Axiom of Replacement applied to α.

$\quad \langle 2 \rangle 5$. γ is a transitive set.

$\quad$ PROOF: If $\mathbf{F}(\delta) < \alpha$ and $\epsilon < \delta$ then $\mathbf{F}(\epsilon) < \alpha$ by $\langle 1 \rangle 2$.

$\quad \langle 2 \rangle 6$. γ is an ordinal.

$\quad$ PROOF: Proposition 6.1.8.

$\quad \langle 2 \rangle 7$. $\mathbf{F}(\gamma) = \alpha$

$\quad \quad \langle 3 \rangle 1$. $\mathbf{F}(\gamma)$ is the least element of $\mathbf{A}$ different from $\mathbf{F}(\delta)$ for all $\delta < \gamma$

$\quad \quad \langle 3 \rangle 2$. $\mathbf{F}(\gamma)$ is the least element of $\mathbf{A}$ different from x for all $x \in \mathbf{A}$ with $x < \alpha$

$\quad \quad \langle 3 \rangle 3$. $\mathbf{F}(\gamma) = \alpha$

$\square$

6.4   Ordinal Operations

Definition 6.4.1 (Ordinal Operation)**.** An *ordinal operation* is a function $\mathbf{On} \to \mathbf{On}$.

Definition 6.4.2 (Continuous)**.** An ordinal operation $\mathbf{T} : \mathbf{On} \to \mathbf{On}$ is *continuous* iff, for every limit ordinal λ, we have $\mathbf{T}(\lambda) = \bigcup_{\alpha < \lambda} \mathbf{T}(\alpha)$.

Definition 6.4.3 (Normal)**.** An ordinal operation is *normal* iff it is continuous and strictly monotone.

Proposition Schema 6.4.4. *For any class* $\mathbf{T}$*, the following is a theorem.*

If $\mathbf{T}$ *is a continuous ordinal operation and* $\forall \gamma . \mathbf{T}(\gamma) < \mathbf{T}(\gamma^+)$*, then* $\mathbf{T}$ *is normal.*

PROOF:

$\langle 1 \rangle 1$. LET: $P[\beta]$ be the property $\forall \gamma < \beta . \mathbf{T}(\gamma) < \mathbf{T}(\beta)$

$\langle 1 \rangle 2$. $P[0]$

$\quad$ PROOF: Vacuous.

$\langle 1 \rangle 3$. For any ordinal γ, if $P[\gamma]$ then $P[\gamma^+]$

$\quad \langle 2 \rangle 1$. ASSUME: $P[\gamma]$

$\quad \langle 2 \rangle 2$. LET: $\delta < \gamma^+$

$\quad \langle 2 \rangle 3$. CASE: $\delta < \gamma$

$\quad$ PROOF: Then $\mathbf{T}(\delta) < \mathbf{T}(\gamma) < \mathbf{T}(\gamma^+)$.

$\quad \langle 2 \rangle 4$. CASE: $\delta = \gamma$

$\quad$ PROOF: Then $\mathbf{T}(\delta) = \mathbf{T}(\gamma) < \mathbf{T}(\gamma^+)$.

$\langle 1 \rangle 4$. For any limit ordinal λ, if $\forall \gamma < \lambda . P[\gamma]$ then $P[\lambda]$.

⟨2⟩1. ASSUME: $\forall \gamma < \lambda.P[\gamma]$
⟨2⟩2. LET: $\delta < \lambda$
⟨2⟩3. $\mathbf{T}(\delta) < \mathbf{T}(\lambda)$
   PROOF:

$$\mathbf{T}(\delta) < \mathbf{T}(\delta^+)$$
$$\leq \bigcup_{\epsilon < \lambda} \mathbf{T}(\epsilon)$$
$$= \mathbf{T}(\lambda)$$

□

Proposition Schema 6.4.5. *For any class* **T**, *the following is a theorem:*
  Assume **T** *is a normal ordinal operation. For every ordinal* α, *we have* $\alpha \leq \mathbf{T}(\alpha)$.

PROOF:
⟨1⟩1. LET: γ be an ordinal.
⟨1⟩2. ASSUME: as induction hypothesis $\forall \delta < \gamma.\mathbf{T}(\delta) \geq \delta$
⟨1⟩3. For all $\delta < \gamma$ we have $\delta < \mathbf{T}(\gamma)$
  PROOF: **T** is strictly monotone.
⟨1⟩4. $\gamma \leq \mathbf{T}(\gamma)$
□

Proposition Schema 6.4.6. *For any class* **T**, *the following is a theorem:*
  Assume **T** *is a normal ordinal operation. For any ordinal* $\beta \geq \mathbf{T}(0)$, *there exists a greatest ordinal* γ *such that* $\mathbf{T}(\gamma) \leq \beta$.

PROOF:
⟨1⟩1. There exists γ such that $\mathbf{T}(\gamma) > \beta$
  ⟨2⟩1. For all γ we have $\mathbf{T}(\gamma) \geq \gamma$
    PROOF: Proposition 6.4.5.
  ⟨2⟩2. $\mathbf{T}(\beta^+) > \beta$
⟨1⟩2. LET: δ be least such that $\mathbf{T}(\delta) > \beta$
⟨1⟩3. δ is a successor ordinal.
  ⟨2⟩1. $\delta \neq 0$
    PROOF: Since $\mathbf{T}(0) \leq \beta$.
  ⟨2⟩2. δ is not a limit ordinal.
    ⟨3⟩1. ASSUME: for a contradiction δ is a limit ordinal.
    ⟨3⟩2. $\beta < \bigcup_{\epsilon < \delta} \mathbf{T}(\epsilon)$
      PROOF: **T** is continuous.
    ⟨3⟩3. There exists $\epsilon < \delta$ such that $\beta < \mathbf{T}(\epsilon)$
    ⟨3⟩4. Q.E.D.
      PROOF: This contradicts the minimality of δ.
⟨1⟩4. LET: $\delta = \gamma^+$
⟨1⟩5. γ is greatest such that $\mathbf{T}(\gamma) \leq \beta$
□

Theorem Schema 6.4.7. *For any class* **T**, *the following is a theorem:*

Assume that $\mathbf{T}$ *is a normal ordinal operation. For any nonempty set of ordinals S, we have*

$$\mathbf{T}(\sup S) = \sup_{\alpha \in S} \mathbf{T}(\alpha) \ .$$

PROOF:

$\langle 1 \rangle 1.$ $\forall \alpha \in S . \mathbf{T}(\alpha) \leq \mathbf{T}(\sup S)$
  PROOF: Since $\mathbf{T}$ is monotone.

$\langle 1 \rangle 2.$ For any ordinal β, if $\forall \alpha \in S . \mathbf{T}(\alpha) \leq \beta$, then $\mathbf{T}(\sup S) \leq \beta$
  $\langle 2 \rangle 1.$ LET: β be an ordinal.
  $\langle 2 \rangle 2.$ LET: $\gamma = \sup S$
  $\langle 2 \rangle 3.$ ASSUME: $\forall \alpha \in S . \mathbf{T}(\alpha) \leq \beta$
  $\langle 2 \rangle 4.$ CASE: γ is 0 or a successor ordinal
    PROOF: Then we must have $\gamma \in S$ so $\mathbf{T}(\gamma) \leq \beta$ from $\langle 2 \rangle 3$.
  $\langle 2 \rangle 5.$ CASE: γ is a limit ordinal
    $\langle 3 \rangle 1.$ $\mathbf{T}(\gamma) = \sup_{\alpha < \gamma} \mathbf{T}(\alpha)$
      PROOF: $\mathbf{T}$ is continuous.
    $\langle 3 \rangle 2.$ ASSUME: for a contradiction $\beta < \mathbf{T}(\gamma)$
    $\langle 3 \rangle 3.$ PICK $\alpha < \gamma$ such that $\beta < \mathbf{T}(\alpha)$
      PROOF: $\langle 3 \rangle 1$, $\langle 3 \rangle 2$
    $\langle 3 \rangle 4.$ PICK $\alpha' \in S$ such that $\alpha < \alpha'$
      PROOF: $\langle 2 \rangle 2$, $\langle 3 \rangle 3$
    $\langle 3 \rangle 5.$ $\beta < \mathbf{T}(\alpha') \leq \beta$
      PROOF: $\mathbf{T}$ is strictly monotone, $\langle 3 \rangle 3$, $\langle 3 \rangle 4$, $\langle 2 \rangle 3$.
    $\langle 3 \rangle 6.$ Q.E.D.
      PROOF: This is a contradiction.

☐

Proposition 6.4.8. *For any classes* $\mathbf{A}$ *and* $\mathbf{T}$, *the following is a theorem:*

Assume $\mathbf{A}$ *is a proper class of ordinals such that, for every set* $S \subseteq \mathbf{A}$, *we have* $\bigcup S \in \mathbf{A}$. *Assume* $\mathbf{T}$ *is the unique order isomorphism* $\mathbf{On} \cong \mathbf{A}$. *Then* $\mathbf{T}$ *is normal.*

PROOF:

$\langle 1 \rangle 1.$ $\mathbf{T}$ is strictly monotone.
  PROOF: Since it is an order isomorphism.

$\langle 1 \rangle 2.$ $\mathbf{T}$ is continuous.
  $\langle 2 \rangle 1.$ LET: λ be a limit ordinal.
  $\langle 2 \rangle 2.$ $\mathbf{T}'(\lambda)$ is the least member of $\mathbf{A}$ that is greater than $\mathbf{T}'(\alpha)$ for all $\alpha < \lambda$
  $\langle 2 \rangle 3.$ $\mathbf{T}'(\lambda) = \sup_{\alpha < \lambda} \mathbf{T}'(\alpha)$

☐

Proposition Schema 6.4.9. *For any class* $\mathbf{T}$, *the following is a theorem:*

If $\mathbf{T}$ *is a normal ordinal operation, then for any limit ordinal* λ, *we have* $\mathbf{T}(\lambda)$ *is a limit ordinal.*

PROOF:

$\langle 1 \rangle 1.$ $\mathbf{T}(\lambda) \neq 0$

PROOF: Since $0 \leq \mathbf{T}(0) < \mathbf{T}(\lambda)$.
$\langle 1 \rangle 2$. $\mathbf{T}(\lambda)$ is not a successor ordinal.
  $\langle 2 \rangle 1$. ASSUME: for a contradiction $\mathbf{T}(\lambda) = \alpha^+$
  $\langle 2 \rangle 2$. $\alpha < \mathbf{T}(\lambda) = \sup_{\beta < \lambda} \mathbf{T}(\beta)$
  $\langle 2 \rangle 3$. PICK $\beta < \lambda$ such that $\alpha < \mathbf{T}(\beta)$
  $\langle 2 \rangle 4$. $\alpha^+ \leq \mathbf{T}(\beta) < \mathbf{T}(\lambda)$
  $\langle 2 \rangle 5$. Q.E.D.
    PROOF: This is a contradiction.
☐

6.5 Ordinal Arithmetic

6.5.1 Addition

Definition 6.5.1. Let A and B be disjoint well-ordered sets. The *concatenation* of A and B is the set $A \cup B$ under the relation:

- if $a, a' \in A$ then $a \leq a'$ iff $a \leq a'$ in A

- if $b, b' \in B$ then $b \leq b'$ iff $b \leq b'$ in B

- if $a \in A$ and $b \in B$ then $a \leq b$ and $b \not\leq a$.

Proposition 6.5.2. *If A and B are disjoint well-ordered sets, then their concatenation is well-ordered.*

PROOF:
$\langle 1 \rangle 1$. $\leq$ is reflexive.
  PROOF: For all $a \in A$ we have $a \leq a$, and for all $b \in B$ we have $b \leq b$.
$\langle 1 \rangle 2$. $\leq$ is antisymmetric.
  $\langle 2 \rangle 1$. ASSUME: $x \leq y \leq x$
  $\langle 2 \rangle 2$. CASE: $x, y \in A$
    PROOF: Then $x = y$ since the order on A is antisymmetric.
  $\langle 2 \rangle 3$. CASE: $x \in A$ and $y \in B$
    PROOF: This is impossible as it would imply $y \not\leq x$.
  $\langle 2 \rangle 4$. CASE: $x \in B$ and $y \in A$
    PROOF: This is impossible as it would imply $x \not\leq y$.
  $\langle 2 \rangle 5$. CASE: $x, y \in B$
    PROOF: Then $x = y$ since the order on B is antisymmetric.
$\langle 1 \rangle 3$. $\leq$ is transitive.
  $\langle 2 \rangle 1$. ASSUME: $x \leq y \leq z$
  $\langle 2 \rangle 2$. CASE: $x, z \in A$
    PROOF: In this case $y \in A$ since $y \leq z$, and so $x \leq z$ since the order on A is transitive.
  $\langle 2 \rangle 3$. CASE: $x \in A$ and $z \in B$
    PROOF: Then $x \leq z$ immediately.
  $\langle 2 \rangle 4$. CASE: $x \in B$ and $z \in A$

PROOF: This is impossible because we have $y \notin A$ since $x \leq y$ and $y \notin B$ since $y \leq z$.

$\langle 2 \rangle 5$. CASE: $x, z \in B$

PROOF: In this case $y \in B$ since $x \leq y$, and so $x \leq z$ since the order on B is transitive.

$\langle 1 \rangle 4$. $\leq$ is total.

$\langle 2 \rangle 1$. LET: $x, y \in A \cup B$

$\langle 2 \rangle 2$. CASE: $x, y \in A$

PROOF: Then $x \leq y$ or $y \leq x$ because the order on A is total.

$\langle 2 \rangle 3$. CASE: $x \in A$ and $y \in B$

PROOF: Then $x \leq y$.

$\langle 2 \rangle 4$. CASE: $x \in B$ and $y \in A$

PROOF: Then $y \leq x$.

$\langle 2 \rangle 5$. CASE: $x, y \in B$

PROOF: Then $x \leq y$ or $y \leq x$ because the order on B is total.

$\langle 1 \rangle 5$. Every nonempty subset of $A \cup B$ has a least element.

$\langle 2 \rangle 1$. LET: S be a nonempty subset of $A \cup B$

$\langle 2 \rangle 2$. CASE: $S \cap A = \emptyset$

PROOF: Then $S \subseteq B$ and so S has a least element.

$\langle 2 \rangle 3$. CASE: $S \cap A \neq \emptyset$

PROOF: The least element of $S \cap A$ is the least element of S.

$\square$

Definition 6.5.3 (Ordinal Addition). Let α and β be ordinal numbers. Then $\alpha + \beta$ is the ordinal number of the concatenation of A and B, where A is any well ordered set with ordinal α and B is any well ordered set with ordinal β.

Theorem 6.5.4 (Associative Law for Addition). *For any ordinals ρ, σ and τ, we have*

$$\rho + (\sigma + \tau) = (\rho + \sigma) + \tau \ .$$

PROOF: Given disjoint well ordered sets A, B and C, the concatenation of A with (the concatenation of B and C) is the same as the concatenation of (the concatenation of A and B) and C. $\square$

Theorem 6.5.5. *For any ordinal ρ we have*

$$\rho + 0 = 0 + \rho = \rho \ .$$

PROOF: For any well ordered set A, the concatenation of A with $\emptyset$ is A, and the concatenation of $\emptyset$ with A is A. $\square$

Theorem 6.5.6. *For any ordinal α we have $\alpha + 1 = \alpha^+$.*

PROOF: Since α^+ is the concatenation of α and $\{\alpha\}$. $\square$

Theorem 6.5.7. *For any ordinal α, the operation that maps β to $\alpha + \beta$ is normal.*

PROOF:

⟨1⟩1. For any limit ordinal λ, we have $\alpha + \lambda = \sup_{\beta < \lambda}(\alpha + \beta)$.

　⟨2⟩1. LET: λ be a limit ordinal.

　⟨2⟩2. $(\{0\} \times \alpha) \cup (\{1\} \times \lambda) = \bigcup_{\beta \in \lambda}((\{0\} \times \alpha) \cup (\{1\} \times \beta))$, where the order on the right hand side is as in Lemma 6.1.15.

　　PROOF:

$$(\{0\} \times \alpha) \cup (\{1\} \times \lambda) = (\{0\} \times \alpha) \cup (\{1\} \times \bigcup_{\beta < \lambda} \beta)$$

$$= (\{0\} \times \alpha) \cup \bigcup_{\beta < \lambda}(\{1\} \times \beta)$$

$$= \bigcup_{\beta < \lambda}((\{0\} \times \alpha) \cup (\{1\} \times \beta))$$

⟨1⟩2. For any ordinal β we have $\alpha + \beta < \alpha + \beta^+$

　PROOF: Since $\alpha + \beta^+ = \alpha + \beta + 1 = (\alpha + \beta)^+$.

☐

Corollary 6.5.7.1. *For any ordinals α, β, γ, we have $\beta < \gamma$ if and only if $\alpha + \beta < \alpha + \gamma$.*

Corollary 6.5.7.2 (Left Cancellation for Addition)**.** *For any ordinals α, β and γ, if $\alpha + \beta = \alpha + \gamma$ then $\beta = \gamma$.*

Theorem 6.5.8. *For any ordinals α, β, γ, if $\beta \leq \gamma$ then $\beta + \alpha \leq \gamma + \alpha$.*

PROOF: Transfinite induction on α. ☐

Theorem 6.5.9 (Subtraction Theorem)**.** *Let α and β be ordinals with $\alpha \leq \beta$. Then there exists a unique ordinal δ such that $\alpha + \delta = \beta$.*

PROOF:

⟨1⟩1. For all ordinals α and β with $\alpha \leq \beta$, there exists δ such that $\alpha + \delta = \beta$

　⟨2⟩1. LET: α and β be ordinals with $\alpha \leq \beta$

　⟨2⟩2. LET: δ be the greatest ordinal such that $\alpha + \delta \leq \beta$

　　PROOF: Proposition 6.4.6.

　⟨2⟩3. $\alpha + \delta = \beta$

　　PROOF: If $\alpha + \delta < \beta$ then $\alpha + \delta + 1 \leq \beta$ contradicting the greatestness of δ.

⟨1⟩2. Q.E.D.

　PROOF: Uniqueness follows from the Left Cancellation Law.

☐

6.5.2　Multiplication

Definition 6.5.10 (Ordinal Multiplication)**.** Let α and β be ordinal numbers. Then $\alpha\beta$ is the ordinal number of $A \times B$ under the lexicographic order, where A is any well ordered set with ordinal α and B is any well ordered set with ordinal β.

　This is well defined by Proposition 5.3.5.

Theorem 6.5.11 (Associative Law)**.** *For any ordinals ρ, σ and τ, we have*

$$\rho(\sigma\tau) = (\rho\sigma)\tau \ .$$

PROOF: Let A, B and C be well ordered sets with ordinals ρ, σ and τ. Then both $\rho(\sigma\tau)$ and $(\rho\sigma)\tau$ are the ordinal of $A \times B \times C$ under
$(a, b, c) \leq (a', b', c') \Leftrightarrow a \leq a' \vee (a = a' \wedge b \leq b') \vee (a = a' \wedge b = b' \wedge c \leq c')$.   □

Theorem 6.5.12 (Left Distributive Law)**.** *For any ordinals ρ, σ and τ, we have*

$$\rho(\sigma + \tau) = \rho\sigma + \rho\tau$$

PROOF: Let A, B and C be well ordered sets with ordinals ρ, σ and τ and with $B \cap C = \emptyset$. Then both $\rho(\sigma + \tau)$ and $\rho\sigma + \rho\tau$ are the ordinal of $A \times (B \cup C)$ under the lexicographic ordering. □

Theorem 6.5.13. *For any ordinal ρ we have $\rho 0 = 0\rho = 0$.*

PROOF: For any well ordered set A we have $A \times \emptyset = \emptyset \times A = \emptyset$. □

Theorem 6.5.14. *For any ordinal ρ we have $\rho 1 = 1\rho = \rho$.*

PROOF: Easy. □

Theorem 6.5.15. *For any ordinals ρ and σ, if $\rho\sigma = 0$ then $\rho = 0$ or $\sigma = 0$.*

PROOF: If $A \times B = \emptyset$ then $A = \emptyset$ or $B = \emptyset$. □

Theorem 6.5.16. *For any non-zero ordinal α, the operation that maps β to $\alpha\beta$ is normal.*

PROOF:
$\langle 1 \rangle 1$. For any limit ordinal λ, we have $\alpha\lambda = \bigcup_{\beta < \lambda} \alpha\beta$
$\quad \langle 2 \rangle 1$. LET: λ be a limit ordinal
$\quad \langle 2 \rangle 2$. $\alpha \times \lambda = \bigcup_{\beta < \lambda} (\alpha \times \beta)$ as well-ordered sets
$\langle 1 \rangle 2$. For any ordinal β we have $\alpha\beta < \alpha\beta^+$
$\quad$ PROOF: $\alpha\beta^+ = \alpha\beta + \alpha > \alpha\beta$
□

Corollary 6.5.16.1. *For any ordinals α, β, γ, if $\alpha \neq 0$ then $\beta < \gamma$ if and only if $\alpha\beta < \alpha\gamma$.*

Corollary 6.5.16.2 (Left Cancellation for Multiplication)**.** *For any ordinals α, β, γ, if $\alpha \neq 0$ and $\alpha\beta = \alpha\gamma$ then $\beta = \gamma$.*

Theorem 6.5.17. *For any ordinals α, β and γ, if $\beta \leq \gamma$ then $\beta\alpha \leq \gamma\alpha$.*

PROOF: Transfinite induction on α. □

Theorem 6.5.18 (Division Theorem)**.** *Let α and δ be ordinal numbers with $\delta \neq 0$. Then there exist unique ordinals β and γ with $\gamma < \delta$ and*

$$\alpha = \delta\beta + \gamma \ .$$

PROOF:

⟨1⟩1. For any ordinal numbers α and δ with $\delta \neq 0$, there exist ordinals β and γ such that $\gamma < \delta$ and $\alpha = \delta\beta + \gamma$

  ⟨2⟩1. LET: α and δ be ordinals with $\delta \neq 0$

  ⟨2⟩2. LET: β be the greatest ordinal such that $\delta\beta \leq \alpha$

    PROOF: Proposition 6.4.6.

  ⟨2⟩3. There exists an ordinal γ such that $\alpha = \delta\beta + \gamma$

    PROOF: Subtraction Theorem

⟨1⟩2. For any ordinals δ, β, β', γ, γ', if $\delta\beta + \gamma = \delta\beta' + \gamma'$ and $\delta \neq 0$ and $\gamma, \gamma' < \delta$ then $\beta = \beta'$ and $\gamma = \gamma'$

  ⟨2⟩1. LET: δ, β, β', γ, γ' be ordinals.

  ⟨2⟩2. ASSUME: $\delta \neq 0$ and $\delta\beta + \gamma = \delta\beta' + \gamma'$

  ⟨2⟩3. $\beta = \beta'$

    ⟨3⟩1. $\beta \not< \beta'$

      PROOF: If $\beta < \beta'$ then

$$\begin{aligned}
\delta\beta' + \gamma' &\geq \delta\beta' \\
&\geq \delta(\beta + 1) \\
&= \delta\beta + \delta \\
&> \delta\beta + \gamma
\end{aligned}$$

    ⟨3⟩2. $\beta' \not< \beta$

      PROOF: Similar.

  ⟨2⟩4. $\gamma = \gamma'$

    PROOF: By Cancellation.

□

6.5.3 Exponentiation

Definition 6.5.19. Given ordinals α and β, define the ordinal α^β as follows:

$$\begin{aligned}
0^\alpha &:= 0 & (\alpha > 0) \\
\alpha^0 &:= 1 & \\
\alpha^{\beta^+} &:= \alpha^\beta \alpha & (\alpha > 0) \\
\alpha^\lambda &:= \sup_{\beta < \lambda} \alpha^\beta & (\alpha > 0, \lambda \text{ a limit ordinal})
\end{aligned}$$

Theorem 6.5.20. *Let α be an ordinal ≥ 2. The operation that maps β to α^β is normal.*

PROOF:

⟨1⟩1. For λ a limit ordinal we have $\alpha^\lambda = \sup_{\beta < \lambda} \alpha^\beta$

  PROOF: By definition.

⟨1⟩2. For any ordinal β we have $\alpha^\beta < \alpha^{\beta^+}$

  PROOF: We have $\alpha^{\beta^+} = \alpha^\beta \alpha > \alpha^\beta$ by Theorem 6.5.16 since $\alpha > 1$ and $\alpha^\beta \neq 0$.

□

Corollary 6.5.20.1. *For any ordinals α, β, γ, if $\alpha \geq 2$ then $\beta < \gamma$ if and only if $\alpha^\beta < \alpha^\gamma$.*

Corollary 6.5.20.2 (Cancellation for Exponentiation)**.** *For any ordinals α, β, γ, if $\alpha \geq 2$ and $\alpha^\beta = \alpha^\gamma$ then $\beta = \gamma$.*

Theorem 6.5.21. *For any ordinals α, β and γ, if $\beta \leq \gamma$ then $\beta^\alpha \leq \gamma^\alpha$.*

PROOF: Transfinite induction on α.

Theorem 6.5.22 (Logarithm Theorem)**.** *Let α and β be ordinal numbers with $\alpha \neq 0$ and $\beta > 1$. Then there exist unique ordinals γ, δ and ρ such that*

$$\alpha = \beta^\gamma \delta + \rho, \qquad 0 \neq \delta < \beta, \qquad \rho < \beta^\gamma \ .$$

PROOF:
$\langle 1 \rangle 1$. For any ordinals α and β with $\alpha \neq 0$ and $\beta > 1$, there exist ordinals γ, δ, ρ such that
$$\alpha = \beta^\gamma \delta + \rho, \qquad 0 \neq \delta < \beta, \qquad \rho < \beta^\gamma \ .$$
  $\langle 2 \rangle 1$. LET: α and β be ordinals with $\alpha \neq 0$ and $\beta > 1$.
  $\langle 2 \rangle 2$. LET: γ be the greatest ordinal such that $\beta^\gamma \leq \alpha$.
    PROOF: Proposition 6.4.6.
  $\langle 2 \rangle 3$. LET: δ and ρ be the unique ordinals with $\rho < \beta^\gamma$ such that $\alpha = \beta^\gamma \delta + \rho$.
    PROOF: By the Division Theorem.
  $\langle 2 \rangle 4$. $\delta \neq 0$
    PROOF: If $\delta = 0$ then $\alpha = \beta^\gamma 0 + \rho = \rho < \beta^\gamma \leq \alpha$ which is a contradiction.
  $\langle 2 \rangle 5$. $\delta < \beta$
    PROOF: If $\beta \leq \delta$ then $\alpha \geq \beta^\gamma \delta \geq \beta^\gamma \beta = \beta^{\gamma+1}$, contradicting the greatestness of γ.
$\langle 1 \rangle 2$. If $\beta^\gamma \delta + \rho = \beta^{\gamma'} \delta' + \rho'$ with $\beta > 1$, $0 \neq \delta < \beta$, $0 \neq \delta' < \beta$, $\rho < \beta^\gamma$ and $\rho' < \beta^{\gamma'}$, then $\gamma = \gamma'$, $\delta = \delta'$ and $\rho = \rho'$.
  $\langle 2 \rangle 1$. LET: $\alpha = \beta^\gamma \delta + \rho = \beta^{\gamma'} \delta' + \rho'$
  $\langle 2 \rangle 2$. $\beta^\gamma \leq \alpha < \beta^{\gamma+1}$
  $\langle 2 \rangle 3$. $\beta^{\gamma'} \leq \alpha < \beta^{\gamma'+1}$
  $\langle 2 \rangle 4$. $\beta^\gamma < \beta^{\gamma'+1}$ and $\beta^{\gamma'} < \beta^{\gamma+1}$
  $\langle 2 \rangle 5$. $\gamma < \gamma' + 1$ and $\gamma' < \gamma + 1$
  $\langle 2 \rangle 6$. $\gamma = \gamma'$
  $\langle 2 \rangle 7$. $\delta = \delta'$ and $\rho = \rho'$
    PROOF: By the Division Theorem.
☐

Theorem 6.5.23. *For any ordinal numbers α, β, γ, we have*

$$\alpha^{\beta+\gamma} = \alpha^\beta \alpha^\gamma \ .$$

PROOF:
$\langle 1 \rangle 1$. LET: $P[\gamma]$ be the property: for any ordinals α and β we have $\alpha^{\beta+\gamma} = \alpha^\beta \alpha^\gamma$
$\langle 1 \rangle 2$. $P[0]$

PROOF:

$$\alpha^{\beta+0} = \alpha^\beta$$
$$= \alpha^\beta 1$$
$$= \alpha^\beta \alpha^0$$

⟨1⟩3. For all γ, if $P[\gamma]$ then $P[\gamma+1]$

PROOF:

$$\alpha^{\beta+\gamma+1} = \alpha^{\beta+\gamma}\alpha$$
$$= \alpha^\beta \alpha^\gamma \alpha \qquad \text{(induction hypothesis)}$$
$$= \alpha^\beta \alpha^{\gamma+1}$$

⟨1⟩4. For any limit ordinal λ, if $\forall \gamma < \lambda.P[\gamma]$ then $P[\lambda]$.

  ⟨2⟩1. LET: λ be a limit ordinal

  ⟨2⟩2. ASSUME: $\forall \gamma < \lambda.P[\gamma]$

  ⟨2⟩3. LET: α and β be any ordinals.

  ⟨2⟩4. CASE: $\alpha = 0$

    PROOF: We have $\alpha^{\beta+\lambda} = \alpha^\beta \alpha^\lambda = 0$.

  ⟨2⟩5. CASE: $\alpha = 1$

    PROOF: We have $\alpha^{\beta+\lambda} = \alpha^\beta \alpha^\lambda = 1$.

  ⟨2⟩6. CASE: $\alpha > 1$

    PROOF:

$$\alpha^{\beta+\lambda} = \alpha^{\sup_{\gamma<\lambda}(\beta+\gamma)}$$
$$= \sup_{\gamma<\lambda} \alpha^{\beta+\gamma} \qquad \text{(Theorem 6.4.7)}$$
$$= \sup_{\gamma<\lambda} \alpha^\beta \alpha^\gamma \qquad (\langle 2\rangle 2)$$
$$= \alpha^\beta \sup_{\gamma<\lambda} \alpha^\gamma \qquad \text{(Theorem 6.4.7)}$$
$$= \alpha^\beta \alpha^\lambda$$

∎

Theorem 6.5.24. *For any ordinal numbers α, β and γ, we have*

$$(\alpha^\beta)^\gamma = \alpha^{\beta\gamma} \ .$$

PROOF:

⟨1⟩1. LET: $P[\gamma]$ be the property: For any ordinals α and β, we have $(\alpha^\beta)^\gamma = \alpha^{\beta\gamma}$

⟨1⟩2. $P[0]$

  PROOF:

$$(\alpha^\beta)^0 = 1$$
$$= \alpha^{\beta 0}$$

⟨1⟩3. $\forall \gamma \in \mathbf{On}.P[\gamma] \Rightarrow P[\gamma+1]$

PROOF:

$$(\alpha^\beta)^{\gamma+1} = (\alpha^\beta)^\gamma \alpha^\beta$$
$$= \alpha^{\beta\gamma} \alpha^\beta$$
$$= \alpha^{\beta\gamma+\beta}$$
$$= \alpha^{\beta(\gamma+1)}$$

⟨1⟩4. For any limit ordinal λ, if $\forall \gamma < \lambda . P[\gamma]$ then $P[\lambda]$.

   ⟨2⟩1. LET: λ be a limit ordinal.

   ⟨2⟩2. ASSUME: $\forall \gamma < \lambda . P[\gamma]$

   ⟨2⟩3. LET: α and β be any ordinals.

   ⟨2⟩4. CASE: $\alpha = 0$ and $\beta = 0$

      PROOF:

$$(0^\beta)^\lambda = 1^\lambda$$
$$= 1$$
$$= 0^0$$
$$= 0^{0\lambda}$$

   ⟨2⟩5. CASE: $\alpha = 0$ and $\beta \neq 0$

      PROOF: $(0^\beta)^\lambda = 0^{\beta\lambda} = 0$.

   ⟨2⟩6. CASE: $\alpha = 1$

      PROOF: $(1^\beta)^\lambda = 1^{\beta\lambda} = 1$

   ⟨2⟩7. CASE: $\alpha > 1$

      PROOF:

$$(\alpha^\beta)^\lambda = \sup_{\gamma<\lambda}(\alpha^\beta)^\gamma$$
$$= \sup_{\gamma<\lambda} \alpha^{\beta\gamma}$$
$$= \alpha^{\sup_{\gamma<\lambda} \beta\gamma}$$
$$= \alpha^{\beta\lambda}$$

⬚

6.6   Sequences

¡

Definition 6.6.1 (Sequence). Given an ordinal α and class **A**, an α-*sequence* in **A** is a function $a : \alpha \to \mathbf{A}$. We write a_β for $a(\beta)$, and $(a_\beta)_{\beta<\alpha}$ for a.

Definition 6.6.2 (Strictly Increasing). A sequence (a_β) of ordinals is *strictly increasing* iff, whenever $\beta < \gamma$, then $a_\beta < a_\gamma$.

Definition 6.6.3 (Subsequence). Let $(a_\beta)_{\beta<\gamma}$ be a sequence in **A**. A *subsequence* of (a_β) is a sequence of the form $(a_{\beta_\xi})_{\xi<\delta}$ where $(\beta_\xi)_{\xi<\delta}$ is a strictly increasing sequence in γ.

Definition 6.6.4 (Convergence). Let $(a_\beta)_{\beta<\gamma}$ be a sequence of ordinals and λ an ordinal. Then (a_β) *converges* to the *limit* λ iff $\lambda = \sup_{\beta<\gamma} a_\beta$.

Lemma 6.6.5. *Let $(a_\beta)_{\beta < \gamma}$ be a sequence of ordinals. Then there is a strictly increasing subsequence $(a_{\beta_\xi})_{\xi < \delta}$ such that $\sup_{\xi < \delta} a_{\beta_\xi} = \sup_{\beta < \gamma} a_\beta$.*

PROOF: Define β_ξ by transfinite recursion as follows. β_ξ is the least β such that $a_\beta > a_{\beta_\zeta}$ for all $\zeta < \xi$ if there is such an a_β; if not, the sequence ends. $\square$

6.7 Strict Supremum

Definition 6.7.1 (Strict Supremum)**.** For any set S of ordinals, define the *strict supremum* of S, ssup S, to be the least ordinal greater than every member of S.

Chapter 7

Cardinal Numbers

7.1 Cardinal Numbers

Definition 7.1.1 (Cardinality). For any set A, the *cardinality* or *cardinal number* $|A|$ of A is the least ordinal equinumerous with A.

Let **Card** be the class of all cardinal numbers.

Proposition 7.1.2. *For any sets A and B, we have $A \approx B$ iff $|A| = |B|$.*

PROOF: Easy. ☐

Definition 7.1.3 (Addition). Given cardinal numbers κ and λ, we define $\kappa + \lambda$ to be $|A \cup B|$ where A and B are disjoint sets of cardinality κ and λ respectively.

We prove this is well-defined.

PROOF:
⟨1⟩1. ASSUME: $A \approx A'$, $B \approx B'$, and $A \cap B = A' \cap B' = \emptyset$
⟨1⟩2. PICK bijections $f : A \approx A'$ and $g : B \approx B'$
⟨1⟩3. The function $A \cup B \to A' \cup B'$ that maps $a \in A$ to $f(a)$ and $b \in B$ to $g(b)$
    is a bijection.
☐

Proposition 7.1.4. *For any cardinal number κ, we have $\kappa + 0 = \kappa$.*

PROOF: Let A and B be disjoint sets of cardinality κ and 0. Then $B = \emptyset$ so $A \cup B = A$ and so $|A \cup B| = \kappa$. ☐

Theorem 7.1.5 (Associative Law for Addition). *For any cardinal numbers κ, λ, μ we have $\kappa + (\lambda + \mu) = (\kappa + \lambda) + \mu$.*

PROOF: Since $A \cup (B \cup C) = (A \cup B) \cup C$. ☐

Proposition 7.1.6. *For any cardinal numbers κ and λ we have $\kappa + \lambda = \lambda + \kappa$.*

PROOF: Since $A \cup B = B \cup A$. ☐

Definition 7.1.7 (Multiplication)**.** For κ and λ cardinal numbers, we define $\kappa\lambda$ to be the cardinal number of $A \times B$, where $|A| = \kappa$ and $|B| = \lambda$.

　　We prove this is well-defined.

PROOF: If $f : A \approx A'$ and $g : B \approx B'$ then the function that maps (a, b) to $(f(a), g(b))$ is a bijection $A \times B \approx A' \times B'$. $\square$

Proposition 7.1.8. *For any cardinal number κ we have $\kappa \cdot 0 = 0$.*

PROOF: Since $A \times \emptyset = \emptyset$. $\square$

Proposition 7.1.9. *For any cardinal number κ we have $\kappa \cdot 1 = \kappa$.*

PROOF: The function that maps (a, e) to a is a bijection $A \times \{e\} \approx A$. $\square$

Theorem 7.1.10 (Distributive Law)**.** *For any cardinal numbers κ, λ and μ, we have $\kappa(\lambda + \mu) = \kappa\lambda + \kappa\mu$.*

PROOF: Since $A \times (B \cup C) = (A \times B) \cup (A \times C)$. $\square$

Theorem 7.1.11 (Associative Law for Multiplication)**.** *For any cardinal numbers κ, λ and μ, we have $\kappa(\lambda\mu) = (\kappa\lambda)\mu$.*

PROOF: Since $A \times (B \times C) \approx (A \times B) \times C$. $\square$

Theorem 7.1.12 (Commutative Law for Multiplication)**.** *For any cardinal numbers κ and λ, we have $\kappa\lambda = \lambda\kappa$.*

PROOF: Since $A \times B \approx B \times A$. $\square$

Theorem 7.1.13. *For any cardinal numbers κ and λ, if $\kappa\lambda = 0$ then $\kappa = 0$ or $\lambda = 0$.*

PROOF: if $A \times B = \emptyset$ then $A = \emptyset$ or $B = \emptyset$. $\square$

Definition 7.1.14 (Exponentiation)**.** Given cardinal numbers κ and λ, we define κ^λ to be $|A^B|$, where $|A| = \kappa$ and $|B| = \lambda$.

　　We prove this is well-defined.

PROOF:If $f : A \approx A'$ and $g : B \approx B'$, then the function that maps $h : B \to A$ to $f \circ h \circ g^{-1}$ is a bijection $A^B \approx A'^{B'}$. $\square$

Proposition 7.1.15. *For any cardinal numbers κ, λ and μ,*

$$\kappa^{\lambda+\mu} = (\kappa^\lambda)^\mu$$

PROOF: The function that maps $f : A \times B \to C$ to $\lambda a \in A.\lambda b \in B.f(a, b)$ is a bijection $A^{B \times C} \approx (A^B)^C$. $\square$

Proposition 7.1.16. *For any cardinal numbers κ, λ and μ,*

$$(\kappa\lambda)^\mu = \kappa^\mu \lambda^\mu \ \ .$$

PROOF: The function $f : A^C \times B^C \to (A \times B)^C$ with $f(g,h)(c) = (g(c), h(c))$ is a bijection. $\square$

Proposition 7.1.17. *For any cardinal numbers κ, λ and μ, we have*

$$\kappa^{\lambda+\mu} = \kappa^\lambda \kappa^\mu \ .$$

PROOF: If $B \cap C = \emptyset$, then $f : A^B \times A^C \to A^{B \cup C}$ given by $f(g,h)(b) = g(b)$ and $f(g,h)(c) = h(c)$ is a bijection. $\square$

Proposition 7.1.18. *For any cardinal number κ, we have $\kappa^0 = 1$.*

PROOF: For any set A, we have $A^\emptyset = \{\emptyset\}$. $\square$

Proposition 7.1.19. *For any cardinal number κ, we have $\kappa^1 = \kappa$.*

PROOF: For any sets A and B, if $B = \{b\}$ then the function $f : A \to A^B$ with $f(a)(b) = a$ is a bijection. $\square$

Proposition 7.1.20. *For any non-zero cardinal number κ we have $0^\kappa = 0$.*

PROOF: If A is nonempty then there is no function $A \to \emptyset$. $\square$

Proposition 7.1.21. *For any set A we have $|\mathcal{P}A| = 2^{|A|}$.*

PROOF: The function $f : \mathcal{P}A \to 2^A$ where $f(X)(a) = 0$ if $a \notin X$ and $f(X)(a) = 1$ if $a \in X$. $\square$

Theorem 7.1.22 (König)**.** *Let I be a set. Let $\{A_i\}_{i \in I}$ and $\{B_i\}_{i \in I}$ be families of sets. Assume that $\forall i \in I.|A_i| < |B_i|$. Then $\left|\bigcup_{i \in I} A_i\right| < \left|\prod_{i \in I} B_i\right|$.*

PROOF:
$\langle 1 \rangle 1$. For all $i \in I$, choose an injection $f_i : A_i \rightarrowtail B_i$
$\langle 1 \rangle 2$. For all $i \in I$, choose $b_i \in B_i - f_i(A_i)$
$\langle 1 \rangle 3$. $\left|\bigcup_{i \in I} A_i\right| \le \left|\prod_{i \in I} B_i\right|$
   $\langle 2 \rangle 1$. Define $g : \bigcup_{i \in I} A_i \to \prod_{i \in I} B_i$ by
$$g(i,a)(j) = \begin{cases} f_i(a) & \text{if } i = j \\ b_j & \text{otherwise} \end{cases}$$
   $\langle 2 \rangle 2$. g is injective.
$\langle 1 \rangle 4$. $\left|\bigcup_{i \in I} A_i\right| \ne \left|\prod_{i \in I} B_i\right|$
   $\langle 2 \rangle 1$. LET: $h : \bigcup_{i \in I} A_i \to \prod_{i \in I} B_i$
      PROVE: h is not surjective.
   $\langle 2 \rangle 2$. For $i \in I$, PICK $c_i \in B_i - \{h(i,a)(i) \mid i \in I\}$
   $\langle 2 \rangle 3$. $c \in \prod_{i \in I} B_i$
   $\langle 2 \rangle 4$. $c \notin \operatorname{ran} h$
$\square$

Corollary 7.1.22.1. *For any cardinal number κ we have $\kappa < 2^\kappa$.*

7.2 Ordering on Cardinal Numbers

Definition 7.2.1. Given cardinal numbers κ and λ, we have $\kappa \leq \lambda$ iff $A \preccurlyeq B$, where $|A| = \kappa$ and $|B| = \lambda$.

PROOF:
$\langle 1 \rangle 1.$ LET: $|A| = \kappa$ and $|B| = \lambda$
$\langle 1 \rangle 2.$ PICK bijections $f : A \approx \kappa$ and $g : B \approx \lambda$
$\langle 1 \rangle 3.$ If $\kappa \leq \lambda$ then $A \preccurlyeq B$
   PROOF: Let $i : \kappa \hookrightarrow \lambda$ be the inclusion. Then $g^{-1} \circ i \circ f$ is an injection $A \to B$.
$\langle 1 \rangle 4.$ If $A \preccurlyeq B$ then $\kappa \leq \lambda$
   $\langle 2 \rangle 1.$ ASSUME: $A \preccurlyeq B$
   $\langle 2 \rangle 2.$ PICK an injection $h : A \rightarrowtail B$
   $\langle 2 \rangle 3.$ $g(h(A)) \subseteq B$ is well-ordered by $\in$
   $\langle 2 \rangle 4.$ LET: γ be the ordinal number of $(g(h(A)), \in)$
   $\langle 2 \rangle 5.$ $\gamma \leq \lambda$
     PROOF: Proposition 6.1.12.
   $\langle 2 \rangle 6.$ $\kappa \leq \gamma$
     PROOF: By the leastness of κ, since A is equinumerous with γ.
   $\langle 2 \rangle 7.$ $\kappa \leq \lambda$
$\square$

Corollary 7.2.1.1. *There is no largest cardinal number.*

Proposition 7.2.2. *For any cardinal numbers κ, λ, μ, if $\kappa \leq \lambda$ then $\kappa + \mu \leq \lambda + \mu$.*

PROOF: If $f : A \to B$ is injective, and $A \cap C = B \cap C = \emptyset$, then the function $A \cup C \to B \cup C$ that maps a to $f(a)$ and maps c to c is an injection. $\square$

Proposition 7.2.3. *For any cardinal numbers κ, λ, μ, if $\kappa \leq \lambda$ then $\kappa\mu \leq \lambda\mu$.*

PROOF: If $f : A \to B$ is injective, then the function $A \times C \to B \times C$ that maps (a, c) to $(f(a), c)$ is injective. $\square$

Proposition 7.2.4. *For any cardinal numbers κ, λ, μ, if $\kappa \leq \lambda$ then $\kappa^\mu \leq \lambda^\mu$.*

PROOF: Given an injection $f : A \to B$, the function that maps $A^C \to B^C$ that maps g to $f \circ g$ is an injection. $\square$

Proposition 7.2.5. *For any cardinal numbers κ, λ, μ, if $\kappa \leq \lambda$ and μ and κ are not both 0, then $\mu^\kappa \leq \mu^\lambda$.*

PROOF:
$\langle 1 \rangle 1.$ LET: A, B and C be sets with A and C not both empty.
$\langle 1 \rangle 2.$ LET: $f : A \to B$ be an injection.
   PROVE: $C^A \preccurlyeq C^B$
$\langle 1 \rangle 3.$ CASE: $C = \emptyset$
  PROOF: Then $A \neq \emptyset$ so $C^A = \emptyset \preccurlyeq C^B$.
$\langle 1 \rangle 4.$ CASE: $C \neq \emptyset$

⟨2⟩1. PICK $c \in C$
⟨2⟩2. LET: $g : C^A \to C^B$ be the function $g(h)(f(a)) = h(a)$, $g(h)(b) = c$ if
     $b \notin f(A)$
⟨2⟩3. g is an injection.
□

Proposition 7.2.6. *Let $\mathcal{A}$ be a set such that $\forall X \in \mathcal{A}.|X| \leq \kappa$. Then*

$$\left| \bigcup \mathcal{A} \right| \leq |\mathcal{A}|\kappa \ .$$

PROOF:
⟨1⟩1. For $X \in \mathcal{A}$, choose a surjection $f_X : \kappa \to X$.
⟨1⟩2. Define $g : \mathcal{A} \times \kappa \to \bigcup \mathcal{A}$ by $g(X, \alpha) = f_X(\alpha)$
⟨1⟩3. g is surjective.
□

Lemma 7.2.7. *The union of a set of cardinal numbers is a cardinal number.*

PROOF:
⟨1⟩1. LET: A be a set of cardinal numbers.
     PROVE: $\bigcup A$ is the smallest ordinal equinumerous with $\bigcup A$
⟨1⟩2. LET: $\alpha < \bigcup A$
     PROVE: $\alpha \not\approx \bigcup A$
⟨1⟩3. PICK $\kappa \in A$ such that $\alpha < \kappa$
⟨1⟩4. $\alpha \prec \kappa$
⟨1⟩5. $\alpha \prec \bigcup A$
□

Chapter 8

Natural Numbers

8.1 Inductive Sets

Definition 8.1.1 (Inductive)**.** A set I is *inductive* iff $0 \in I$ and $\forall x \in I.x^+ \in I$.

Definition 8.1.2 (Natural Number)**.** A *natural number* is a set that belongs to every inductive set.

Theorem 8.1.3. *The class $\mathbb{N}$ of natural numbers is a set.*

Proof:
⟨1⟩1. Pick an inductive set I.
   Proof: Axiom of Infinity.
⟨1⟩2. $\mathbb{N} \subseteq I$
☐

Theorem 8.1.4. $\mathbb{N}$ *is inductive, and is a subset of every other inductive set.*

Proof:
⟨1⟩1. $\mathbb{N}$ is inductive.
  ⟨2⟩1. $0 \in \mathbb{N}$
    Proof: Since 0 is a member of every inductive set.
  ⟨2⟩2. $\forall n \in \mathbb{N}.n^+ \in \mathbb{N}$
    ⟨3⟩1. Let: $n \in \mathbb{N}$
    ⟨3⟩2. Let: I be any inductive set.
       Prove: $n^+ \in I$
    ⟨3⟩3. $n \in I$
      Proof: ⟨3⟩1, ⟨3⟩2
    ⟨3⟩4. $n^+ \in I$
      Proof: ⟨3⟩2, ⟨3⟩3
⟨1⟩2. $\mathbb{N}$ is a subset of every inductive set.
  Proof: Immediate from definitions.
☐

Corollary 8.1.4.1 (Induction Principle for $\mathbb{N}$)**.** *Any inductive subset of $\mathbb{N}$ coincides with $\mathbb{N}$.*

Theorem 8.1.5. *Every natural number except 0 is the successor of some natural number.*

PROOF: Trivially by induction. $\square$

Proposition 8.1.6. *Every natural number is an ordinal.*

PROOF: By induction. $\square$

Proposition 8.1.7. $\mathbb{N}$ *is a transitive set.*

PROOF:
$\langle 1 \rangle 1$. $0 \subseteq \mathbb{N}$
$\langle 1 \rangle 2$. $\forall n \in \mathbb{N}.n \subseteq \mathbb{N} \Rightarrow n^+ \subseteq \mathbb{N}$
$\langle 1 \rangle 3$. $\forall n \in \mathbb{N}.n \subseteq \mathbb{N}$
   PROOF: From $\langle 1 \rangle 1$ and $\langle 1 \rangle 2$ by induction.
$\square$

Corollary 8.1.7.1. $\mathbb{N}$ *is an ordinal.*

Definition 8.1.8. We define $\omega = \mathbb{N}$.

Proposition 8.1.9 (Dependent Choice)**.** *Let A be a nonempty set and R a relation on A such that $\forall x \in A.\exists y \in A.(y, x) \in R$. Then there exists a function $f : \mathbb{N} \to A$ such that $\forall n \in \mathbb{N}.(f(n+1), f(n)) \in R$.*

PROOF:
$\langle 1 \rangle 1$. PICK a choice function F for A.
$\langle 1 \rangle 2$. PICK $a \in A$
$\langle 1 \rangle 3$. Define $f : \mathbb{N} \to A$ by $f(0) = a$ and $f(n+1) = F(\{y \in A \mid (y, f(n)) \in R\})$.
$\square$

Theorem Schema 8.1.10. *For any classes $\mathbf{A}$ and $\mathbf{R}$, the following is a theorem:*

   Assume $\mathbf{R}$ is a relation on $\mathbf{A}$ and, for all $a \in \mathbf{A}$, the class $\{x \in \mathbf{A} \mid x\mathbf{R}a\}$ is a set. Then $\mathbf{R}$ is well founded if and only if there does not exist a function $f : \mathbb{N} \to \mathbf{A}$ such that $\forall n \in \mathbb{N}.f(n+1)\mathbf{R}f(n)$.

PROOF:
$\langle 1 \rangle 1$. If there exists a function $f : \mathbb{N} \to \mathbf{A}$ such that $\forall n \in \mathbb{N}.f(n+1)\mathbf{R}f(n)$ then
      $\mathbf{R}$ is not well founded.
   PROOF: $f(\mathbb{N})$ is a nonempty subset of $\mathbf{A}$ with no $\mathbf{R}$-minimal element.
$\langle 1 \rangle 2$. If $\mathbf{R}$ is not well founded then there exists a function $f : \mathbb{N} \to \mathbf{A}$ such that
      $\forall n \in \mathbb{N}.f(n+1)\mathbf{R}f(n)$.
   $\langle 2 \rangle 1$. ASSUME: $\mathbf{R}$ is not well founded.
   $\langle 2 \rangle 2$. PICK a nonempty subset $B \subseteq \mathbf{A}$ that has no $\mathbf{R}$-minimal element.
   $\langle 2 \rangle 3$. $\forall x \in B.\exists y \in B.y\mathbf{R}x$

$\langle 2 \rangle 4$. Choose a function $g : B \to B$ such that $\forall x \in B.g(x)\mathbf{R}x$
$\langle 2 \rangle 5$. PICK $b \in B$
$\langle 2 \rangle 6$. Define $f : \mathbb{N} \to \mathbf{A}$ recursively by $f(0) = b$ and $\forall n \in \mathbb{N}.f(n{+}1) = g(f(n))$
$\langle 2 \rangle 7$. $\forall n \in \mathbb{N}.f(n + 1)\mathbf{R}f(n)$
☐

8.2   Cardinality

Definition 8.2.1 (Finite). A set is *finite* iff it is equinumerous to some natural number; otherwise it is *infinite*.

Theorem 8.2.2 (Pigeonhole Principle). *No natural number is equinumerous to a proper subset of itself.*

PROOF:
$\langle 1 \rangle 1$. LET: $P(n)$ be the property: any one-to-one function $n \to n$ is surjective.
$\langle 1 \rangle 2$. $P(0)$
   PROOF: The only function $0 \to 0$ is injective.
$\langle 1 \rangle 3$. For every natural number n, if $P(n)$ then $P(n + 1)$.
   $\langle 2 \rangle 1$. ASSUME: $P(n)$
   $\langle 2 \rangle 2$. LET: f be a one-to-one function $n + 1 \to n + 1$
   $\langle 2 \rangle 3$. $f \upharpoonright n$ is a one-to-one function $n \to n + 1$
   $\langle 2 \rangle 4$. CASE: $n \notin ranf$
      $\langle 3 \rangle 1$. $f \upharpoonright n : n \to n$
      $\langle 3 \rangle 2$. $\mathrm{ran}(f \upharpoonright n) = n$
      $\langle 3 \rangle 3$. $f(n) = n$
         PROOF: $\langle 2 \rangle 1$.
      $\langle 3 \rangle 4$. $\mathrm{ran}\, f = n + 1$
   $\langle 2 \rangle 5$. CASE: $n \in \mathrm{ran}\, f$
      $\langle 3 \rangle 1$. PICK $p \in n$ such that $f(p) = n$
      $\langle 3 \rangle 2$. LET: $\hat{f} : n \to n$ be the function
$$\hat{f}(p) = f(n)$$
$$\hat{f}(x) = f(x) \qquad\qquad (x \neq p)$$
      $\langle 3 \rangle 3$. $\hat{f}$ is one-to-one
      $\langle 3 \rangle 4$. $\mathrm{ran}\, \hat{f} = n$
         PROOF: $\langle 2 \rangle 1$
      $\langle 3 \rangle 5$. $\mathrm{ran}\, f = n + 1$
$\langle 1 \rangle 4$. For every natural number n, $P(n)$.
☐

Corollary 8.2.2.1. *No finite set is equinumerous to a proper subset of itself.*

Corollary 8.2.2.2. *Every natural number is a cardinal number.*

PROOF: For any natural number n, we have that n is the least ordinal such that $n \approx n$. ☐

Corollary 8.2.2.3. $\mathbb{N}$ *is a cardinal number.*

Corollary 8.2.2.4. $\mathbb{N}$ *is infinite.*

PROOF: The function that maps n to $n+1$ is a bijection between $\mathbb{N}$ and $\mathbb{N}-\{0\}$. $\square$

Corollary 8.2.2.5. *If C is a proper subset of a natural number n, then there exists $m < n$ such that $C \approx m$.*

PROOF: By Proposition 6.1.12. $\square$

Corollary 8.2.2.6. *Any subset of a finite set is finite.*

Proposition 8.2.3. *For any natural numbers m and n we have $m+n$ (cardinal addition) is a natural number.*

PROOF: Induction on n. $\square$

Corollary 8.2.3.1. *The union of two finite sets is finite.*

Corollary 8.2.3.2. *The union of a finite set of finite sets is finite.*

PROOF: By induction on the number of elements. $\square$

Proposition 8.2.4. *For natural numbers m and n, the cardinal sum $m + n$ is equal to the ordinal sum $m + n$.*

PROOF: Induction on n. $\square$

Proposition 8.2.5. *For any natural numbers m and n, we have mn (cardinal multiplication) is a natural number.*

Corollary 8.2.5.1. *If A and B are finite sets then $A \times B$ is finite.*

Proposition 8.2.6. *For natural numbers m and n, the cardinal product mn is equal to the ordinal product mn.*

PROOF: Induction on n. $\square$

Proposition 8.2.7. *For any natural numbers m and n we have m^n (cardinal exponentiation) is a natural number.*

PROOF: Induction on n.

Corollary 8.2.7.1. *If A and B are finite sets then A^B are finite.*

Proposition 8.2.8. *For natural numbers m and n, the cardinal exponentiation m^n and the ordinal exponentiation m^n agree.*

PROOF: Induction on n. $\square$

Proposition 8.2.9. $\mathbb{N}^2 \approx \mathbb{N}$

PROOF: The function $J : \mathbb{N}^2 \to \mathbb{N}$ defined by $J(m, n) = ((m+n)^2 + 3m + n)/2$ is a bijection. $\square$

Proposition 8.2.10. *For any infinite cardinal κ we have $\aleph_0 \leq \kappa$.*

PROOF:
$\langle 1 \rangle 1.$ LET: A be an infinite set.
     PROVE: $\mathbb{N} \preccurlyeq A$
$\langle 1 \rangle 2.$ PICK a choice function F for A.
$\langle 1 \rangle 3.$ Define $h : \mathbb{N} \to \{X \in \mathcal{P}A \mid X \text{ is finite}\}$ by
$$h(0) = \emptyset$$
$$h(n+1) = h(n) \cup \{F(A - \{h(m) \mid m < n\})\}$$
$\langle 1 \rangle 4.$ Define $g : \mathbb{N} \to A$ by $g(n) = F(A - \{h(m) \mid m < n\})$
$\langle 1 \rangle 5.$ g is injective.
    PROOF: If $m < n$ then $g(m) \neq g(n)$.
$\square$

Theorem Schema 8.2.11 (König's Lemma)**.** *For any classes $\mathbf{A}$ and $\mathbf{R}$, the following is a theorem:*

    Assume $\mathbf{R}$ is a well founded relation on $\mathbf{A}$ such that, for all $y \in \mathbf{A}$, the class $\{x \in \mathbf{A} \mid x\mathbf{R}y\}$ is a finite set. Let $\mathbf{R}^t$ be the transitive closure of $\mathbf{R}$. Then, for all $y \in \mathbf{A}$, the class $\{x \in \mathbf{A} \mid x\mathbf{R}^t y\}$ is a finite set.

PROOF:
$\langle 1 \rangle 1.$ LET: $y \in \mathbf{A}$
$\langle 1 \rangle 2.$ ASSUME: as transfinite induction hypothesis $\forall x\mathbf{R}y.\{z \in \mathbf{A} \mid z\mathbf{R}^t x\}$ is a
             finite set.
$\langle 1 \rangle 3.$ $\{x \mid x\mathbf{R}^t y\} = \bigcup_{x\mathbf{R}y}(\{x\} \cup \{z \mid z\mathbf{R}^t x\}$
$\langle 1 \rangle 4.$ $\{x \mid x\mathbf{R}^t y\}$ is finite.
    PROOF: Corollary 8.2.3.2.
$\square$

8.3 Countable Sets

Definition 8.3.1 (Countable)**.** A set A is *countable* iff $|A| \leq \aleph_0$.

Theorem 8.3.2. *The union of a countable set of countable sets is countable.*

PROOF: Proposition 7.2.6. $\square$

8.4 Arithmetic

Definition 8.4.1 (Even)**.** A natural number n is *even* iff there exists $m \in \mathbb{N}$ such that $n = 2m$.

Definition 8.4.2 (Odd)**.** A natural number n is *odd* iff there exists $p \in \mathbb{N}$ such that $n = 2p + 1$.

Proposition 8.4.3. *Every natural number is either even or odd.*

PROOF:

⟨1⟩1. 0 is even.
  PROOF: $0 = 2 \times 0$.
⟨1⟩2. For any natural number n, if n is either even or odd then n^+ is either even
      or odd.
  PROOF:
  ⟨2⟩1. LET: $n \in \mathbb{N}$
  ⟨2⟩2. If n is even then n^+ is odd.
    PROOF: If $n = 2p$ then $n^+ = 2p + 1$.
  ⟨2⟩3. If n is odd then n^+ is even.
    PROOF: If $n = 2p + 1$ then $n^+ = 2(p + 1)$.
☐

Proposition 8.4.4. *No natural number is both even and odd.*

PROOF:

⟨1⟩1. 0 is not odd.
  PROOF: For any p we have $2p + 1 = (2p)^+ \neq 0$.
⟨1⟩2. For any natural number n, if n is not both even and odd, then n^+ is not
      both even and odd.
  ⟨2⟩1. LET: n be a natural number.
  ⟨2⟩2. If n^+ is even then n is odd.
    ⟨3⟩1. ASSUME: n^+ is even.
    ⟨3⟩2. PICK p such that $n^+ = 2p$
    ⟨3⟩3. $p \neq 0$
      PROOF: Since $n^+ \neq 0$.
    ⟨3⟩4. PICK q such that $p = q^+$
      PROOF: Theorem 8.1.5.
    ⟨3⟩5. $n^+ = 2q + 2$
      PROOF: ⟨3⟩2, ⟨3⟩4.
    ⟨3⟩6. $n = 2q + 1$
      PROOF: Proposition 6.2.7, ⟨3⟩5
    ⟨3⟩7. n is odd.
  ⟨2⟩3. If n^+ is odd then n is even.
    ⟨3⟩1. ASSUME: n^+ is odd.
    ⟨3⟩2. PICK p such that $n^+ = 2p + 1$
    ⟨3⟩3. $n = 2p$
      PROOF: Proposition 6.2.7, ⟨3⟩2
    ⟨3⟩4. n is even.
☐

Proposition 8.4.5. *Let m, n, p, q be natural numbres. Assume $m + n = p + q$.*
Then $m < p$ if and only if $q < n$.

PROOF:

⟨1⟩1. If $m < p$ then $q < n$.

PROOF: If $m < p$ and $n \le q$ then $m + n < p + q$.
⟨1⟩2. If $q < n$ then $m < p$.
  PROOF: Similar.
☐

Proposition 8.4.6. *Let m, n, p and q be natural numbers. Assume $n < m$ and $q < p$. Then*

$$mq + np < mp + nq \ .$$

PROOF:
⟨1⟩1. PICK positive natural numbers a and b such that $m = n + a$ and $p = q + b$.
⟨1⟩2. $mp + nq > mq + np$
  PROOF:

$$
\begin{aligned}
mp + nq &= (n + a)(q + b) + nq \\
&= 2nq + nb + aq + ab \\
mq + np &= (n + a)q + n(q + b) \\
&= 2nq + aq + nb \\
\therefore mp + nq &= mq + np + ab \\
&> mq + np
\end{aligned}
$$

☐

8.5   Sequences

Definition 8.5.1 (Sequence)**.** Let A be a set. A *finite sequence* in A is a function $a : n \to A$ for some natural number n; we write it as $(a(0), a(1), \ldots, a(n - 1))$. An *(infinite) sequence* in A is a function $\mathbb{N} \to A$.
  We write A^* for the set of all finite sequences in A.

Proposition 8.5.2. *If A is countable then A^* is countable.*

PROOF: For any n, the set A^n is countable, and A^* is equinumerous with $\bigcup_n A^n$.
☐

8.6   Transitive Closure of a Set

Proposition 8.6.1. *For any set A, there exists a unique transitive set C such that:*

- *$A \subseteq C$*

- *For any transitive set X, if $A \subseteq X$ then $C \subseteq X$*

PROOF:
⟨1⟩1. Define a function $F : \mathbb{N} \to \mathbf{V}$ by
$$F(0) = A$$

$$F(n + 1) = A \cup \bigcup (F(0) \cup \cdots \cup F(n))$$

⟨1⟩2. For all $n \in \mathbb{N}$ and $a \in F(n)$ we have $a \subseteq F(n+1)$
   PROOF: $a \in F(0) \cup \cdots \cup F(n)$ so $a \subseteq \bigcup(F(0) \cup \cdots \cup F(n)) \subseteq F(n+1)$.
⟨1⟩3. LET: $C = \bigcup_{n \in \mathbb{N}} F(n)$
⟨1⟩4. C is transitive.
   ⟨2⟩1. LET: $x \in y \in C$
   ⟨2⟩2. PICK $n \in \mathbb{N}$ such that $y \in F(n)$
   ⟨2⟩3. $y \subseteq F(n+1)$
      PROOF: ⟨1⟩2
   ⟨2⟩4. $x \in F(n+1)$
   ⟨2⟩5. $x \in C$
⟨1⟩5. $A \subseteq C$
   PROOF: Since $F(0) = A$.
⟨1⟩6. For any transitive set X, if $A \subseteq X$ then $C \subseteq X$
   ⟨2⟩1. LET: X be a transitive set
   ⟨2⟩2. ASSUME: $A \subseteq X$
   ⟨2⟩3. For all $n \in \mathbb{N}$ we have $F(n) \subseteq X$.
      ⟨3⟩1. $F(0) \subseteq X$
         PROOF: ⟨2⟩2
      ⟨3⟩2. For all $n \in \mathbb{N}$, if $F(n) \subseteq X$, then $F(n+1) \subseteq X$.
         ⟨4⟩1. LET: $n \in \mathbb{N}$
         ⟨4⟩2. ASSUME: $\forall m < n.F(m) \subseteq X$
         ⟨4⟩3. $F(0) \cup \cdots \cup F(n) \subseteq X$
         ⟨4⟩4. $\bigcup(F(0) \cup \cdots \cup F(n)) \subseteq X$
            PROOF: Since X is transitive.
         ⟨4⟩5. $F(n+1) \subseteq X$
   ⟨2⟩4. $C \subseteq X$
⟨1⟩7. Let D be a transitive set such that $A \subseteq D$ and, for any transitive set X,
   if $A \subseteq X$ then $D \subseteq X$. Then $D = C$.
   PROOF: We have $C \subseteq D$ and $D \subseteq C$.
☐

8.7   The Veblen Fixed Point Theorem

Theorem Schema 8.7.1 (Veblen Fixed Point Theorem). *For any class* **T***, the following is a theorem:*

Assume **T** *is a normal ordinal operation. For every ordinal β, there exists $\gamma \geq \beta$ such that* $\mathbf{T}(\gamma) = \gamma$.

PROOF:
⟨1⟩1. LET: β be an ordinal.
⟨1⟩2. ASSUME: w.l.o.g. $\beta < \mathbf{T}(\beta)$
   PROOF: We have $\beta \leq \mathbf{T}(\beta)$ by Proposition 6.4.5, and if $\beta = \mathbf{T}(\beta)$ we take
   $\gamma := \beta$.

⟨1⟩3. Define $f : \mathbb{N} \to \mathbf{On}$ by recursion thus:
$$f(0) = \beta$$
$$f(n^+) = \mathbf{T}(f(n))$$

⟨1⟩4. LET: $\gamma = \sup_{n \in \mathbb{N}} f(n)$

⟨1⟩5. $\beta \leq \gamma$

  PROOF: Since $\beta = f(0)$.

⟨1⟩6. $\mathbf{T}(\gamma) = \gamma$

  ⟨2⟩1. $\mathbf{T}(\gamma) \leq \gamma$

    PROOF:
$$\begin{aligned}
\mathbf{T}(\gamma) &= \sup_{n \in \mathbb{N}} \mathbf{T}(f(n)) && \text{(Theorem 6.4.7)} \\
&= \sup_{n \in \mathbb{N}} f(n^+) && (\langle 1 \rangle 3) \\
&\leq \sup_{n \in \mathbb{N}} f(n) \\
&= \gamma
\end{aligned}$$

  ⟨2⟩2. $\gamma \leq \mathbf{T}(\gamma)$

    PROOF:Proposition 6.4.5.

◻

Definition 8.7.2 (Derived Operation). Let $\mathbf{T}$ be a normal ordinal operation. The *derived* operation $\mathbf{T}' : \mathbf{On} \to \mathbf{V}$ is the unique order isomorphism between $\mathbf{On}$ and the fixed points of $\mathbf{T}$.

Proposition Schema 8.7.3. *For any class* $\mathbf{T}$, *the following is a theorem:*
  If $\mathbf{T}$ *is a normal ordinal operation, then the derived operation is normal.*

PROOF:

⟨1⟩1. For any set S of fixed points of $\mathbf{T}$, we have $\bigcup S$ is a fixed point of $\mathbf{T}$

  ⟨2⟩1. LET: S be a set of fixed points of $\mathbf{T}$.

  ⟨2⟩2. $\mathbf{T}(\sup S) = \sup S$

    PROOF:
$$\begin{aligned}
\mathbf{T}(\sup S) &= \sup_{\alpha \in S} \mathbf{T}(\alpha) && \text{(Theorem 6.4.7)} \\
&= \sup_{\alpha \in S} \alpha && (\langle 2 \rangle 1) \\
&= \sup S
\end{aligned}$$

⟨1⟩2. Q.E.D.

  PROOF: Proposition 6.4.8.

◻

8.8 Cantor Normal Form

Theorem 8.8.1. *For any ordinal* α, *there exist a unique sequence of nonzero natural numbers* $(n_1, \ldots, n_k)$ *and sequence of ordinals* $(\gamma_1, \ldots, \gamma_k)$ *such that*

$$\gamma_k < \gamma_{k-1} < \cdots < \gamma_1$$

and

$$\alpha = \omega^{\gamma_1} n_1 + \omega^{\gamma_2} n_2 + \cdots + \omega^{\gamma_k} n_k \ .$$

PROOF:

⟨1⟩1. For any ordinal α, there exist a sequence of nonzero natural numbers $(n_1, \ldots, n_k)$ and sequence of ordinals $(\gamma_1, \ldots, \gamma_k)$ such that
$$\gamma_k < \gamma_{k-1} < \cdots < \gamma_1$$
and
$$\alpha = \omega^{\gamma_1} n_1 + \omega^{\gamma_2} n_2 + \cdots + \omega^{\gamma_k} n_k \ .$$

⟨2⟩1. LET: α be an ordinal

⟨2⟩2. ASSUME: as an induction hypothesis that, for all $\beta < \alpha$, the theorem holds.

⟨2⟩3. ASSUME: w.l.o.g. $\alpha \neq 0$

⟨2⟩4. LET: γ_1, n_1, ρ_1 be the unique ordinals such that $0 \neq n_1 < \omega$, $\rho_1 < \omega^{\gamma_1}$, and $\alpha = \omega^{\gamma_1} n_1 + \rho_1$

⟨2⟩5. LET: $(\gamma_2, \ldots, \gamma_k)$ and $(n_2, \ldots, n_k)$ be sequences such that $\gamma_k < \gamma_{k-1} < \cdots < \gamma_2$ and $\rho_1 = \omega^{\gamma_2} n_2 + \cdots + \omega^{\gamma_k} n_k$

⟨2⟩6. $\gamma_2 < \gamma_1$

　　PROOF: Since $\omega^{\gamma_2} \leq \rho_1 < \omega^{\gamma_1}$

⟨1⟩2. If
$$\gamma_k < \gamma_{k-1} < \cdots < \gamma_1$$
$$\gamma_k' < \gamma_{k-1}' < \cdots < \gamma_1'$$
and
$$\omega^{\gamma_1} n_1 + \omega^{\gamma_2} n_2 + \cdots + \omega^{\gamma_k} n_k = \omega^{\gamma_1'} n_1' + \omega^{\gamma_2'} n_2' + \cdots + \omega^{\gamma_k'} n_k'$$
then $\gamma_i = \gamma_i'$ for all i and $n_i = n_i'$ for all i

　PROOF: Prove by induction on i using the Logarithm Theorem.

∎

Definition 8.8.2 (Cantor Normal Form)**.** For any ordinal α, the *Cantor normal form* of α is the expression $\alpha = \omega^{\gamma_1} n_1 + \cdots + \omega^{\gamma_k} n_k$ such that $n_1, \ldots, n_k$ are nonzero natural numbers and $\gamma_k < \gamma_{k-1} < \cdots < \gamma_1$.

Chapter 9

The Cumulative Hierarchy

Definition 9.0.1 (Cumulative Hierarchy). Define the function $V : \mathbf{On} \to \mathbf{V}$ by transfinite recursion thus:

$$V_\alpha = \bigcup_{\beta < \alpha} \mathcal{P}V_\beta$$

Proposition 9.0.2. *For all $\alpha \in \mathbf{On}$, V_α is a transitive set.*

PROOF:
$\langle 1 \rangle 1.$ LET: $\alpha \in \mathbf{On}$
$\langle 1 \rangle 2.$ ASSUME: as transfinite induction hypothesis $\forall \beta < \alpha . V_\beta$ is a transitive set.
$\langle 1 \rangle 3.$ For all $\beta < \alpha$, $\mathcal{P}V_\beta$ is a transitive set.
  PROOF: Proposition 1.6.4.
$\langle 1 \rangle 4.$ V_α is a transitive set.
  PROOF: Proposition 1.6.3.
□

Proposition 9.0.3. *For any ordinals α and β, if $\beta < \alpha$ then $V_\beta \subseteq V_\alpha$.*

PROOF: Since $V_\beta \in \mathcal{P}V_\beta \subseteq V_\alpha$ and V_α is a transitive set. □

Theorem 9.0.4.

  1. $V_0 = \emptyset$

  2. $\forall \alpha \in \mathbf{On}.V_{\alpha^+} = \mathcal{P}V_\alpha$

  3. *For any limit ordinal λ, $V_\lambda = \bigcup_{\alpha < \lambda} V_\alpha$.*

PROOF:
$\langle 1 \rangle 1.$ $V_0 = \emptyset$
  PROOF: Immediate from definition.
$\langle 1 \rangle 2.$ $\forall \alpha \in \mathbf{On}.V_{\alpha^+} = \mathcal{P}V_\alpha$
  PROOF:

$\langle 2 \rangle 1$. LET: $\alpha \in \mathbf{On}$

$\langle 2 \rangle 2$. For all $\beta < \alpha$ we have $\mathcal{P}V_\beta \subseteq \mathcal{P}V_\alpha$

   PROOF: Propositions 1.5.8 and 9.0.3.

$\langle 2 \rangle 3$. $V_{\alpha^+} = \mathcal{P}V_\alpha$

$$
\begin{aligned}
V_{\alpha^+} &= \bigcup_{\beta < \alpha^+} \mathcal{P}V_\beta \\
&= \bigcup_{\beta < \alpha} \mathcal{P}V_\beta \cup \mathcal{P}V_\alpha \\
&\quad \mathcal{P}V_\alpha
\end{aligned}
$$
<div align="right">□</div>

$\langle 1 \rangle 3$. For any limit ordinal λ, $V_\lambda = \bigcup_{\alpha < \lambda} V_\alpha$

   PROOF:

$\langle 2 \rangle 1$. $V_\lambda \subseteq \bigcup_{\alpha < \lambda} V_\alpha$

     PROOF:

$$
\begin{aligned}
V_\lambda &= \bigcup_{\alpha < \lambda} \mathcal{P}V_\alpha \\
&= \bigcup_{\alpha < \lambda} V_{\alpha^+} \qquad\qquad (\langle 1 \rangle 2) \\
&\subseteq \bigcup_{\alpha < \lambda} V_\alpha
\end{aligned}
$$

$\langle 2 \rangle 2$. $\bigcup_{\alpha < \lambda} V_\alpha \subseteq V_\lambda$

     PROOF: Proposition 9.0.3.

□

Proposition 9.0.5. *For every set A, there exists an ordinal α such that $A \in V_\alpha$.*

PROOF:

$\langle 1 \rangle 1$. Let us say a set A is *grounded* iff there exists an ordinal α such that $A \in V_\alpha$.

$\langle 1 \rangle 2$. For any set A, if every element of A is grounded, then A is grounded.

   $\langle 2 \rangle 1$. LET: A be a set.

   $\langle 2 \rangle 2$. $S = \{\alpha \mid \exists a \in A . \alpha \text{ is the least ordinal such that } a \in V_\alpha\}$

     PROOF: S is a set by an Axiom of Replacement.

   $\langle 2 \rangle 3$. LET: $\beta = \sup S$

   $\langle 2 \rangle 4$. $A \subseteq V_\beta$

     $\langle 3 \rangle 1$. LET: $a \in A$

     $\langle 3 \rangle 2$. LET: α be the least ordinal such that $a \in V_\beta$

     $\langle 3 \rangle 3$. $\alpha \in S$

     $\langle 3 \rangle 4$. $\alpha \leq \beta$

     $\langle 3 \rangle 5$. $a \in V_\beta$

   $\langle 2 \rangle 5$. $A \in V_{\beta^+}$

$\langle 1 \rangle 3$. ASSUME: for a contradiction there exists an ungrounded set.

$\langle 1 \rangle 4$. PICK a transitive set B that has an ungrounded member.

   PROOF: Pick a transitive set c, and take B to be the transitive closure of $\{c\}$.

$\langle 1 \rangle 5$. LET: $A = \{x \in B \mid x \text{ is ungrounded}\}$

$\langle 1 \rangle 6$. PICK $m \in A$ such that $m \cap A = \emptyset$
  PROOF: Axiom of Regularity.
$\langle 1 \rangle 7$. Every member of m is grounded.
  $\langle 2 \rangle 1$. ASSUME: for a contradiction $x \in m$ is ungrounded.
  $\langle 2 \rangle 2$. $x \in B$
    PROOF: Since B is transitive ($\langle 1 \rangle 4$).
  $\langle 2 \rangle 3$. $x \in A$
    PROOF: $\langle 1 \rangle 5$
  $\langle 2 \rangle 4$. Q.E.D.
    PROOF: This contradicts $\langle 1 \rangle 6$.
$\langle 1 \rangle 8$. m is grounded.
  PROOF: $\langle 1 \rangle 2$
$\langle 1 \rangle 9$. Q.E.D.
  PROOF: This contradicts $\langle 1 \rangle 6$.
☐

Definition 9.0.6 (Rank). The *rank* of a set A is the least ordinal α such that $A \in V_{\alpha^+}$.

Proposition 9.0.7. *For any set A we have*

$$\operatorname{rank} A = \bigcup_{a \in A} (\operatorname{rank} a)^+$$

PROOF:
$\langle 1 \rangle 1$. LET: $\alpha = \bigcup_{a \in A} (\operatorname{rank} a)^+$
$\langle 1 \rangle 2$. $A \subseteq V_\alpha$
  $\langle 2 \rangle 1$. LET: $a \in A$
  $\langle 2 \rangle 2$. $a \in V_{(\operatorname{rank} a)^+}$
  $\langle 2 \rangle 3$. $a \in V_\alpha$
$\langle 1 \rangle 3$. $A \in V_{\alpha^+}$
$\langle 1 \rangle 4$. If $A \subseteq V_\beta$ then $\alpha \leq \beta$
  $\langle 2 \rangle 1$. ASSUME: $A \subseteq V_\beta$
  $\langle 2 \rangle 2$. For all $a \in A$ we have $(\operatorname{rank} a)^+ \leq \beta$
    PROOF: Since $a \in V_\beta$.
  $\langle 2 \rangle 3$. $\alpha \leq \beta$
☐

Corollary 9.0.7.1. *For any sets a and b, if $a \in b$ then $\operatorname{rank} a < \operatorname{rank} b$.*

Proposition 9.0.8. *For any ordinal number α we have $\operatorname{rank} \alpha = \alpha$.*

PROOF:
$\langle 1 \rangle 1$. LET: α be an ordinal.
$\langle 1 \rangle 2$. ASSUME: as transfinite induction hypothesis $\forall \beta < \alpha. \operatorname{rank} \beta = \beta$
$\langle 1 \rangle 3$. $\operatorname{rank} \alpha = \bigcup_{\beta < \alpha} \beta^+$

PROOF:

$$\operatorname{rank} \alpha = \bigcup_{\beta < \alpha} (\operatorname{rank} \beta)^+$$

$$= \bigcup_{\beta < \alpha} \beta^+$$

$\langle 1 \rangle 4.$ $\bigcup_{\beta < \alpha} \beta^+ \leq \alpha$
    PROOF: Since for all $\beta < \alpha$ we have $\beta^+ \leq \alpha$.
$\langle 1 \rangle 5.$ $\alpha \leq \bigcup_{\beta < \alpha} \beta^+$
    $\langle 2 \rangle 1.$ LET: $\gamma = \bigcup_{\beta < \alpha} \beta^+$
    $\langle 2 \rangle 2.$ ASSUME: for a contradiction $\gamma < \alpha$
    $\langle 2 \rangle 3.$ $\gamma^+ \leq \bigcup_{\beta < \alpha} \beta^+ = \gamma$
    $\langle 2 \rangle 4.$ Q.E.D.
      PROOF: This is a contradiction.

⬚

Definition 9.0.9 (Hereditarily Finite)**.** A set is *hereditarily finite* iff it is in V_ω.

Chapter 10

Models of Set Theory

Definition 10.0.1 (Relativization). Let σ be a sentence in the language of set theory and $\mathbf{M}$ a class. The *relativization* of σ to $\mathbf{M}$ is the sentence $\sigma^{\mathbf{M}}$ formed by replacing every quantifier $\forall x$ with $\forall x \in \mathbf{M}$, and $\exists x$ with $\exists x \in \mathbf{M}$.

We write '$\mathbf{M}$ is a model of σ' for the sentence $\sigma^{\mathbf{M}}$.

Theorem Schema 10.0.2. *For any class* $\mathbf{M}$, *the following is a theorem:*

If $\mathbf{M}$ *is a transitive class, then* $\mathbf{M}$ *is a model of the Axiom of Extensionality.*

PROOF:
$\langle 1 \rangle 1$. ASSUME: $\mathbf{M}$ is a transitive class.
   PROVE: $\forall x, y \in \mathbf{M}(\forall z \in \mathbf{M}(z \in x \Leftrightarrow z \in y) \Rightarrow x = y)$
$\langle 1 \rangle 2$. LET: $x, y \in \mathbf{M}$
$\langle 1 \rangle 3$. ASSUME: $\forall z \in \mathbf{M}(z \in x \Leftrightarrow z \in y)$
$\langle 1 \rangle 4$. $\forall z(z \in x \Leftrightarrow z \in y)$
   PROOF: Since $z \in x \Rightarrow z \in \mathbf{M}$ and $z \in y \Rightarrow z \in \mathbf{M}$ by $\langle 1 \rangle 1$.
$\langle 1 \rangle 5$. $x = y$
$\square$

Theorem 10.0.3. *If* α *is a non-zero ordinal then* V_α *is a model of the statement: The empty class is a set.*

PROOF:
$\langle 1 \rangle 1$. LET: $\alpha \neq 0$
   PROVE: $\exists x \in V_\alpha . \forall y \in V_\alpha . y \notin x$
$\langle 1 \rangle 2$. $\emptyset \in V_\alpha$
$\langle 1 \rangle 3$. $\forall y \in V_\alpha . y \notin \emptyset$
$\square$

Theorem 10.0.4. *For any limit ordinal* λ, *we have* V_λ *is a model of the statement: for any sets* a *and* b, *the class* $\{a, b\}$ *is a set.*

PROOF:
$\langle 1 \rangle 1$. LET: λ be a limit ordinal.

PROVE: $\forall a, b \in V_\lambda.\exists c \in V_\lambda.\forall x \in V_\lambda(x \in c \Leftrightarrow x = a \vee x = b)$

$\langle 1 \rangle 2.$ LET: $a, b \in V_\lambda$

$\langle 1 \rangle 3.$ PICK $\alpha, \beta < \lambda$ such that $a \in V_\alpha$ and $b \in V_\beta$

$\langle 1 \rangle 4.$ ASSUME: w.l.o.g. $\alpha \leq \beta$

$\langle 1 \rangle 5.$ $a, b \in V_\beta$

$\langle 1 \rangle 6.$ $\{a, b\} \in V_{\beta+1}$

$\langle 1 \rangle 7.$ $\{a, b\} \in V_\lambda$

$\langle 1 \rangle 8.$ $\forall x \in V_\lambda(x \in \{a, b\} \Leftrightarrow x = a \vee x = b)$

□

Theorem 10.0.5. *For any ordinal α, we have V_α is a model of the Union Axiom.*

PROOF:

$\langle 1 \rangle 1.$ LET: α be an ordinal.
PROVE: $\forall a \in V_\alpha.\exists b \in V_\alpha.\forall x \in V_\alpha(x \in b \Leftrightarrow \exists y \in V_\alpha(x \in y \wedge y \in a))$

$\langle 1 \rangle 2.$ LET: $a \in V_\alpha$

$\langle 1 \rangle 3.$ PICK $\beta < \alpha$ such that $a \subseteq V_\beta$

$\langle 1 \rangle 4.$ $\bigcup a \subseteq V_\beta$
 PROOF: V_β is a transitive set.

$\langle 1 \rangle 5.$ $\bigcup a \in V_\alpha$

$\langle 1 \rangle 6.$ $\forall x \in V_\alpha(x \in \bigcup a \Leftrightarrow \exists y \in V_\alpha(x \in y \wedge y \in a))$
 PROOF: V_α is a transitive set.

□

Theorem 10.0.6. *For any limit ordinal λ, we have V_λ is a model of the Power Set Axiom.*

PROOF:

$\langle 1 \rangle 1.$ LET: λ be a limit ordinal.
PROVE: $\forall a \in V_\lambda.\exists b \in V_\lambda.\forall x \in V_\lambda(x \in b \Leftrightarrow \forall y \in V_\lambda(y \in x \Rightarrow y \in a))$

$\langle 1 \rangle 2.$ LET: $a \in V_\lambda$

$\langle 1 \rangle 3.$ PICK $\alpha < \lambda$ such that $a \in V_\alpha$

$\langle 1 \rangle 4.$ $\mathcal{P}a \in V_{\alpha+1}$

$\langle 1 \rangle 5.$ $\mathcal{P}a \in V_\lambda$

$\langle 1 \rangle 6.$ $\forall x \in V_\lambda(x \in \mathcal{P}a \Leftrightarrow \forall y \in V_\lambda(y \in x \Rightarrow y \in a))$

□

Theorem Schema 10.0.7. *For any property $P[x, y_1, \ldots, y_n]$, the following is a theorem:*

   For any ordinal α, the set V_α is a model of the statement: for any sets a_1, $\ldots, a_n$, B, the class $\{x \in B \mid P[x, a_1, \ldots, a_n]\}$ is a set.

PROOF:

$\langle 1 \rangle 1.$ LET: α be an ordinal.

$\langle 1 \rangle 2.$ LET: $a_1, \ldots, a_n, B \in V_\alpha$

$\langle 1 \rangle 3.$ LET: $C = \{x \in B \mid P[x, a_1, \ldots, a_n]^{V_\alpha}\}$

$\langle 1 \rangle 4.$ $C \in V_\alpha$

$\langle 1\rangle 5.$ $\forall x \in V_\alpha(x \in C \Leftrightarrow x \in B \land P[x, a_1, \dots, a_n]^{V_\alpha})$
□

Theorem 10.0.8. *For any ordinal $\alpha > \omega$, we have: V_α is a model of the Axiom of Infinity.*

PROOF:
$\langle 1\rangle 1.$ LET: $\alpha > \omega$
$\langle 1\rangle 2.$ $\mathbb{N} \in V_\alpha$
$\langle 1\rangle 3.$ $\exists e \in V_\alpha(e \in \mathbb{N} \land \forall x \in V_\alpha.x \notin e)$
$\langle 1\rangle 4.$ $\forall x \in V_\alpha(x \in \mathbb{N} \Rightarrow \exists y \in V_\alpha \forall z \in V_\alpha(z \in y \Leftrightarrow z \in x \lor z = x))$
□

Theorem 10.0.9. *For any ordinal α, we have V_α is a model of the Axiom of Choice.*

PROOF:
$\langle 1\rangle 1.$ LET: α be an ordinal.
$\langle 1\rangle 2.$ LET: $A \in V_\alpha$
$\langle 1\rangle 3.$ ASSUME: $\forall x \in V_\alpha(x \in A \Rightarrow \exists y \in V_\alpha.y \in A)$
$\langle 1\rangle 4.$ ASSUME: $\forall x, y, z \in V_\alpha(x \in A \land y \in A \land z \in x \land z \in y \Rightarrow x = y)$
$\langle 1\rangle 5.$ A is a set of pairwise disjoint nonempty sets.
$\langle 1\rangle 6.$ PICK c such that, for all $x \in A$, $x \cap c = \emptyset$
$\langle 1\rangle 7.$ $c \cap \bigcup A \in V_\alpha$
$\langle 1\rangle 8.$ $\forall x \in V_\alpha(x \in A \Rightarrow \exists y \in V_\alpha \forall z \in V_\alpha(z \in x \land z \in c \cap \bigcup A \Leftrightarrow z = y))$
□

Theorem 10.0.10. *For any ordinal α, we have V_α is a model of the Axiom of Regularity.*

PROOF:
$\langle 1\rangle 1.$ LET: α be an ordinal.
$\langle 1\rangle 2.$ LET: $A \in V_\alpha$
$\langle 1\rangle 3.$ ASSUME: $\exists x \in V_\alpha.x \in A$
$\langle 1\rangle 4.$ PICK $m \in A$ of least rank.
$\langle 1\rangle 5.$ $m \in V_\alpha$
$\langle 1\rangle 6.$ $\neg \exists x \in V_\alpha(x \in m \land x \in A)$
□

Theorem Schema 10.0.11. *For any axiom α of Zermelo set theory, the following is a theorem:*
    For any limit ordinal $\lambda > \omega$, we have V_λ is a model of α.

PROOF: Theorems 10.0.2, 10.0.3, 10.0.4, 10.0.5, 10.0.6, 10.0.7, 10.0.8, 10.0.9, 10.0.10. □

Corollary Schema 10.0.11.1. *for any axiom α of Zermelo set theory, the following is a theorem:*
    $V_{\omega 2}$ is a model of α.

Lemma 10.0.12. *There exists a well-ordered structure in $V_{\omega 2}$ whose ordinal is not in $V_{\omega 2}$.*

PROOF: Take the well-ordered set $\mathbb{N} \times \{0, 1\}$ whose ordinal is $\omega 2$. $\square$

Corollary Schema 10.0.12.1. *There exists an instance α of the Axiom Schema of Replacement such that the following is a theorem:*

   $V_{\omega 2}$ is not a model of α.

Chapter 11

Infinite Cardinals

11.1 Arithmetic of Infinite Cardinals

Proposition 11.1.1. *For any infinite cardinal κ we have $\kappa\kappa = \kappa$.*

PROOF:

$\langle 1 \rangle 1$. PICK a set B with $|B| = \kappa$

$\langle 1 \rangle 2$. LET: $\mathcal{H} = \{f \mid f = \emptyset \vee \exists A \subseteq B.(A$ is infinite $\wedge f : A \times A \approx A\}$

$\langle 1 \rangle 3$. For any chain $\mathcal{C} \subseteq \mathcal{H}$ we have $\bigcup \mathcal{C} \in \mathcal{H}$

  $\langle 2 \rangle 1$. LET: $\mathcal{C} \subseteq \mathcal{H}$ be a chain.

  $\langle 2 \rangle 2$. ASSUME: w.l.o.g. $\mathcal{C}$ has a nonempty element.

  $\langle 2 \rangle 3$. $\bigcup \mathcal{C}$ is a function.

    $\langle 3 \rangle 1$. ASSUME: $(x, y), (x, z) \in \bigcup \mathcal{C}$

    $\langle 3 \rangle 2$. PICK $f, g \in \mathcal{C}$ such that $f(x) = y$ and $g(x) = z$

    $\langle 3 \rangle 3$. ASSUME: w.l.o.g. $f \subseteq g$

    $\langle 3 \rangle 4$. $y = z$

  $\langle 2 \rangle 4$. $\bigcup \mathcal{C}$ is injective.

    PROOF: Similar.

  $\langle 2 \rangle 5$. LET: $A = \operatorname{ran} \bigcup \mathcal{C}$

  $\langle 2 \rangle 6$. A is infinite.

    $\langle 3 \rangle 1$. PICK a nonzero $f \in \mathcal{C}$

    $\langle 3 \rangle 2$. LET: A' be the infinite subset of B such that $f : A'^2 \approx A'$

    $\langle 3 \rangle 3$. $A' \subseteq A$

  $\langle 2 \rangle 7$. $\operatorname{dom} \bigcup \mathcal{C} = A^2$

    $\langle 3 \rangle 1$. LET: $x, y \in A$

    $\langle 3 \rangle 2$. PICK $f, g \in \mathcal{C}$ such that $x \in \operatorname{ran} f$ and $y \in \operatorname{ran} g$

    $\langle 3 \rangle 3$. ASSUME: w.l.o.g. $f \subseteq g$

    $\langle 3 \rangle 4$. LET: A' be the infinite subset of B such that $g : A'^2 \approx A'$

    $\langle 3 \rangle 5$. $x, y \in A'$

    $\langle 3 \rangle 6$. $(x, y) \in \operatorname{dom} g$

    $\langle 3 \rangle 7$. $(x, y) \in \operatorname{dom} \bigcup \mathcal{C}$

  $\langle 2 \rangle 8$. $\bigcup \mathcal{C} \in \mathcal{H}$

$\langle 1 \rangle 4$. PICK a maximal $f_0 \in \mathcal{H}$

$\langle 1 \rangle 5$. $f_0 \neq \emptyset$

   $\langle 2 \rangle 1$. PICK a countably infinite subset A of B.

     PROOF: Proposition 8.2.10.

   $\langle 2 \rangle 2$. PICK a bijection $f : A^2 \approx A$

     PROOF: Proposition 8.2.9.

   $\langle 2 \rangle 3$. $\emptyset \subseteq f \in \mathcal{H}$

   $\langle 2 \rangle 4$. $\emptyset$ is not maximal in $\mathcal{H}$

$\langle 1 \rangle 6$. LET: A_0 be the infinite subset of B such that $f_0 : A_0^2 \approx A_0$

$\langle 1 \rangle 7$. LET: $\lambda = |A_0|$

$\langle 1 \rangle 8$. λ is infinite.

$\langle 1 \rangle 9$. $\lambda^2 = \lambda$

$\langle 1 \rangle 10$. $\lambda = \kappa$

   $\langle 2 \rangle 1$. ASSUME: for a contradiction $\lambda < \kappa$

   $\langle 2 \rangle 2$. $\lambda \leq |B - A_0|$

   $\langle 2 \rangle 3$. PICK a subset $D \subseteq B - A_0$ with $|D| = \lambda$

   $\langle 2 \rangle 4$. $(A_0 \cup D)^2 = A_0^2 \cup (A_0 \times D) \cup (D \times A_0) \cup D^2$

   $\langle 2 \rangle 5$. LET: $C = (A_0 \times D) \cup (D \times A_0) \cup D^2$

   $\langle 2 \rangle 6$. $|C| = \lambda$

     PROOF:

$$
\begin{aligned}
|(A_0 \times D) \cup (D \times A_0) \cup D^2| &= \lambda^2 + \lambda^2 + \lambda^2 \\
&= \lambda + \lambda + \lambda && (\langle 1 \rangle 9) \\
&= 3\lambda \\
&\leq \lambda \cdot \lambda \\
&= \lambda && (\langle 1 \rangle 9)
\end{aligned}
$$

   $\langle 2 \rangle 7$. PICK a bijection $g : C \approx D$

   $\langle 2 \rangle 8$. $f_0 \cup g : (A_0 \cup D)^2 \approx A_0 \cup D$

   $\langle 2 \rangle 9$. Q.E.D.

     PROOF: This contradicts the maximality of f_0.

   □

Theorem 11.1.2 (Absorption Law of Cardinal Arithmetic). *Let κ and λ be nonzero cardinal numbers such that at least one is infinite. Then*

$$
\kappa + \lambda = \kappa\lambda = \max(\kappa, \lambda)
$$

PROOF:

$\langle 1 \rangle 1$. ASSUME: w.l.o.g. $\lambda \leq \kappa$

$\langle 1 \rangle 2$. $\kappa + \lambda = \kappa\lambda = \kappa$

PROOF:

$$\begin{aligned}
\kappa &\leq \kappa + \lambda \\
&\leq \kappa + \kappa \\
&= 2\kappa \\
&\leq \kappa\lambda \\
&\leq \kappa\kappa \\
&= \kappa \qquad\qquad\qquad \text{(Proposition 11.1.1)}
\end{aligned}$$

□

11.2 Alephs

Definition 11.2.1 (Aleph)**.** Let $\aleph$ be the unique order isomorphism between **On** and the class of infinite cardinals.

Proposition 11.2.2. *The operation $\aleph$ is normal.*

PROOF: Proposition 6.4.8 and Lemma 7.2.7. □

Definition 11.2.3 (Continuum Hypothesis)**.** The *continuum hypothesis* is the statement that $\aleph_1 = 2^{\aleph_0}$.

Definition 11.2.4 (Generalised Continuum Hypothesis)**.** The *generalised continuum hypothesis* is the statement that, for all α, $\aleph_{\alpha^+} = 2^{\aleph_\alpha}$.

11.3 Beths

Definition 11.3.1 (Beth)**.** Define the operation $\beth : \mathbf{On} \to \mathbf{Card}$ by transfinite recursion as follows:

$$\begin{aligned}
\beth_0 &:= \aleph_0 \\
\beth_{\alpha^+} &:= 2^{\beth_\alpha} \\
\beth_\lambda &:= \bigcup_{\alpha < \lambda} \beth_\alpha \qquad\qquad (\lambda \text{ a limit ordinal})
\end{aligned}$$

Proposition 11.3.2. $\beth$ *is a normal operation.*

PROOF: It is continuous by definition, and $\beth_\alpha < \beth_{\alpha^+}$ by Cantor's Theorem. □

Proposition 11.3.3. *The continuum hypothesis is equivalent to the statement* $\beth_1 = \aleph_1$.
  The generalised continuum hypothesis is equivalent to the statement $\beth = \alpha$.

PROOF: Immediate from definitions. □

Lemma 11.3.4. *For any ordinal number α, we have $|V_{\omega+\alpha}| = \beth_\alpha$.*

PROOF:

$\langle 1 \rangle 1.$ $|V_\omega| = \beth_0$
　　PROOF: Since V_ω is the union of $\aleph_0$ finite sets of increasing size.
$\langle 1 \rangle 2.$ For any ordinal α, if $|V_{\omega+\alpha}| = \beth_\alpha$ then $|V_{\omega+\alpha+1}| = \beth_{\alpha+1}$
　　PROOF: Since $V_{\omega+\alpha+1} = \mathcal{P}V_{\omega+\alpha}$.
$\langle 1 \rangle 3.$ For any limit ordinal λ, if $\forall \alpha < \lambda.|V_{\omega+\alpha}| = \beth_\alpha$ then $|V_{\omega+\lambda}| = \beth_\lambda$.
　　PROOF:

$$
\begin{aligned}
|V_{\omega+\lambda}| &= \left| \bigcup_{\alpha<\lambda} V_{\omega+\alpha} \right| \\
&= \sup_{\alpha<\lambda} |V_{\omega+\alpha}| \\
&= \sup_{\alpha<\lambda} \beth_\alpha \\
&= \beth_\lambda
\end{aligned}
$$

$\square$

11.4　Cofinality

Definition 11.4.1 (Cofinal)**.** Let λ be a limit ordinal and S a set of ordinals smaller than λ. Then S is *cofinal* in λ if and only if $\lambda = \sup S$.

Definition 11.4.2 (Cofinality)**.** For any ordinal α, define the *cofinality* of α, cf α, as follows:

- cf $0 = 0$

- For any ordinal α, cf $\alpha^+ = 1$

- For any limit ordinal λ, cf λ is the smallest cardinal such that there exists a set S of ordinals cofinal in λ with $|S| = $ cf λ.

Definition 11.4.3 (Regular)**.** A cardinal κ is *regular* iff cf $\kappa = \kappa$; otherwise it is *singular*.

Proposition 11.4.4. $\aleph_0$ *is regular.*

PROOF: $\aleph_0$ is not the supremum of $< \aleph_0$ smaller ordinals, because a finite union of finite ordinals is finite. $\square$

Proposition 11.4.5. *For every ordinal α, $\aleph_{\alpha+1}$ is regular.*

PROOF: If S is a set of ordinals with $|S| < \aleph_{\alpha+1}$ and $\forall \beta \in S.\beta < \aleph_{\alpha+1}$, then we have $|S| \leq \aleph_\alpha$ and $\forall \beta \in S.\beta \leq \aleph_\alpha$, hence

$$
\begin{aligned}
\left| \bigcup S \right| &\leq \aleph_\alpha^2 &&\text{(Proposition 7.2.6)} \\
&= \aleph_\alpha &&\text{(Proposition 11.1.1)} \square
\end{aligned}
$$

Proposition Schema 11.4.6. *For any class* **T***, the following is a theorem.*
　　Assume **T** $:$ **On** $\to$ **On** *is a normal operation. For any limit ordinal λ we have* cf **T**$(\lambda) = $ cf λ.

PROOF:

$\langle 1 \rangle 1$. cf $\mathbf{T}(\lambda) \leq$ cf λ

$\quad \langle 2 \rangle 1$. PICK a set S of ordinals $< \lambda$ with $|S| =$ cf λ and sup $S = \lambda$

$\quad \langle 2 \rangle 2$. $\mathbf{T}(\lambda) = \sup_{\alpha \in S} \mathbf{T}(\alpha)$

$\quad\quad$ PROOF: Theorem 6.4.7.

$\langle 1 \rangle 2$. cf $\lambda \leq$ cf $\mathbf{T}(\lambda)$

$\quad \langle 2 \rangle 1$. PICK a set A of ordinals $< \mathbf{T}(\lambda)$ such that $|A| =$ cf $\mathbf{T}(\lambda)$ and sup $A = \mathbf{T}(\lambda)$

$\quad \langle 2 \rangle 2$. LET: $B = \{\gamma < \lambda \mid \exists \alpha \in A . |\alpha| = \mathbf{T}(\gamma)\}$

$\quad \langle 2 \rangle 3$. $|B| \leq |A| =$ cf $\mathbf{T}(\lambda)$

$\quad\quad$ PROVE: sup $B = \lambda$

$\quad \langle 2 \rangle 4$. $\forall \alpha \in A . |\alpha| \leq \mathbf{T}(\sup B)$

$\quad \langle 2 \rangle 5$. $\forall \alpha \in A . \alpha < \mathbf{T}(\sup B + 1)$

$\quad \langle 2 \rangle 6$. $\aleph_\lambda = \sup A \leq \mathbf{T}(\sup B + 1)$

$\quad \langle 2 \rangle 7$. $\lambda \leq \sup B + 1$

$\quad \langle 2 \rangle 8$. $\lambda \leq \sup B$

$\quad\quad$ PROOF: λ is a limit ordinal.

$\quad \langle 2 \rangle 9$. sup $B = \lambda$

$\square$

Corollary 11.4.6.1. $\aleph_\omega$ *is singular.*

PROOF: cf $\aleph_\omega =$ cf $\aleph_0 = \aleph_0$. $\square$

Corollary 11.4.6.2. *The operation* cf *is not strictly monotone or continuous.*

PROOF: cf $\aleph_\omega <$ cf $\aleph_1$ $\square$

Definition 11.4.7 (Weakly Inaccessible)**.** A cardinal is *weakly inaccessible* iff it is $\aleph_\lambda$ for some limit ordinal λ and regular.

Lemma 11.4.8. *Let λ be a limit ordinal. Then there exists a strictly increasing* cf λ-*sequence that converges to λ.*

PROOF:

$\langle 1 \rangle 1$. PICK a set S of ordinals $< \lambda$ with $|S| =$ cf λ and sup $S = \lambda$

$\langle 1 \rangle 2$. PICK a bijection $a :$ cf $\lambda \approx S$

$\langle 1 \rangle 3$. PICK a strictly increasing subsequence $(b_\delta)_{\delta < \beta}$ of a that converges to λ.

$\quad$ PROOF: Lemma 6.6.5.

$\langle 1 \rangle 4$. $\beta =$ cf λ

$\quad$ PROOF: By minimiality of cf λ.

$\square$

Corollary 11.4.8.1. *Let λ be a limit ordinal. Then* cf λ *is the least ordinal such that there exists a strictly increasing* cf λ-*sequence that converges to λ.*

Proposition 11.4.9. *For any ordinal λ,* cf λ *is a regular cardinal.*

PROOF:

⟨1⟩1. LET: λ be an ordinal.
⟨1⟩2. ASSUME: w.l.o.g. λ is a limit ordinal.
⟨1⟩3. PICK a strictly increasing sequence $(a_\alpha)_{\alpha < \mathrm{cf}\,\lambda}$ that converges to λ.
⟨1⟩4. LET: S be a set of ordinals $< \mathrm{cf}\,\lambda$ such that $|S| = \mathrm{cf}\,\mathrm{cf}\,\lambda$ and $\sup S = \mathrm{cf}\,\lambda$.
⟨1⟩5. LET: $a(S) = \{a_\alpha \mid \alpha \in S\}$
⟨1⟩6. $a(S)$ is cofinal in λ.
  ⟨2⟩1. LET: $\beta < \lambda$
  ⟨2⟩2. PICK $\gamma < \mathrm{cf}\,\lambda$ such that $\beta < a_\gamma$
  ⟨2⟩3. PICK $\delta \in S$ such that $\gamma < \delta$
  ⟨2⟩4. $a_\delta \in a(S)$ and $\beta < a_\gamma < a_\delta$
⟨1⟩7. $\mathrm{cf}\,\lambda \leq \mathrm{cf}\,\mathrm{cf}\,\lambda$
  PROOF: Since $a(S)$ is a set of ordinals $< \lambda$ with $|a(S)| = \mathrm{cf}\,\mathrm{cf}\,\lambda$ and $\sup a(S) = \lambda$.
⟨1⟩8. $\mathrm{cf}\,\mathrm{cf}\,\lambda = \mathrm{cf}\,\lambda$
☐

Theorem 11.4.10. *Let λ be an infinite cardinal. Then $\mathrm{cf}\,\lambda$ is the least cardinal such that λ can be partitioned into $\mathrm{cf}\,\lambda$ sets, each of cardinality $< \lambda$.*

PROOF:
⟨1⟩1. λ can be partitioned into $\mathrm{cf}\,\lambda$ sets, each of cardinality $< \lambda$
  ⟨2⟩1. PICK a strictly increasing sequence of ordinlas $(a_\alpha)_{\alpha < \mathrm{cf}\,\lambda}$ that converges to λ
  ⟨2⟩2. $\{\{\beta \mid a_\alpha \leq \beta < a_{\alpha+1}\} \mid \alpha < \mathrm{cf}\,\lambda\}$ is a partition of λ into $\mathrm{cf}\,\lambda$ sets, each of cardinality $< \lambda$
⟨1⟩2. If λ can be partitioned into κ sets, each of cardinality $< \lambda$, then $\mathrm{cf}\,\lambda \leq \kappa$.
  ⟨2⟩1. LET: $\mathcal{A}$ be a partition of λ into sets of cardinality $< \lambda$
  ⟨2⟩2. LET: $\kappa = |P|$
  ⟨2⟩3. PICK a bijection $A : \kappa \approx P$
  ⟨2⟩4. $\lambda = \bigcup_{\xi < \kappa} A(\xi)$
  ⟨2⟩5. For all $\xi < \kappa$ we have $|A(\xi)| < \lambda$
  ⟨2⟩6. LET: $\mu = \sup_{\xi < \kappa} |A(\xi)|$
  ⟨2⟩7. $\mu \leq \lambda$
  ⟨2⟩8. For all $\xi < \kappa$ we have $|A(\xi)| \leq \mu$
  ⟨2⟩9. $\lambda \leq \mu\kappa$
    PROOF: Proposition 7.2.6.
  ⟨2⟩10. ASSUME: w.l.o.g. $\kappa < \lambda$
    PROOF: If $\lambda \leq \kappa$ then $\mathrm{cf}\,\lambda \leq \kappa$ since $\mathrm{cf}\,\lambda \leq \lambda$.
  ⟨2⟩11. $\lambda = \mu$
    PROOF:

$$\begin{aligned}
\lambda &\leq \mu\kappa & (\langle 2\rangle 9)\\
&\leq \lambda\lambda & (\langle 2\rangle 7, \langle 2\rangle 10)\\
&= \lambda & (\text{Proposition 11.1.1})
\end{aligned}$$

  ⟨2⟩12. $\{|A(\xi)| \mid \xi < \kappa\}$ is a set of $\leq \kappa$ ordinals all $< \lambda$ whose supremum is λ
  ⟨2⟩13. $\mathrm{cf}\,\lambda \leq \kappa$
☐

Theorem 11.4.11 (König). *For any infinite cardinal κ we have $\kappa < \operatorname{cf} 2^\kappa$.*

PROOF:
$\langle 1 \rangle 1$. ASSUME: for a contradiction $\operatorname{cf} 2^\kappa \leq \kappa$
$\langle 1 \rangle 2$. LET: $S = 2^\kappa$
$\langle 1 \rangle 3$. PICK a partition $\{A_\xi \mid \xi < \kappa\}$ of S^κ with $\forall \xi < \kappa. |A_\xi| < 2^\kappa$.
 PROOF: Theorem 11.4.10.
$\langle 1 \rangle 4$. $\forall \xi < \kappa. \{g(\xi) \mid g \in A_\xi\} \subsetneq S$
 PROOF: We do not have equality because $|\{g(\xi) \mid g \in A_\xi\}| \leq |A_\xi| < 2^\kappa$.
$\langle 1 \rangle 5$. For all $\xi < \kappa$, choose $s_\xi \in S - \{g(\xi) \mid g \in A_\xi\}$
$\langle 1 \rangle 6$. $s \in S^\kappa$
$\langle 1 \rangle 7$. For all $\xi < \kappa$ we have $s \notin A_\xi$
 PROOF: Since for all $\xi < \kappa$ and $g \in A_\xi$ we have $s_\xi(\xi) \neq g(\xi)$.
$\langle 1 \rangle 8$. Q.E.D.
 PROOF: This contradicts $\langle 1 \rangle 3$.
$\square$

Corollary 11.4.11.1.
$$2^{\aleph_0} \neq \aleph_\omega$$

Proposition 11.4.12. *For any ordinal α, we have $\operatorname{cf} \alpha$ is the least cardinal such that α is the strict supremum of $\operatorname{cf} \alpha$ smaller ordinals.*

PROOF:
$\langle 1 \rangle 1$. CASE: $\alpha = 0$
 PROOF: Since $0 = \operatorname{ssup} \emptyset$.
$\langle 1 \rangle 2$. CASE: $\alpha = \beta^+$
 PROOF: Since $\beta^+ = \operatorname{ssup}\{\beta\}$.
$\langle 1 \rangle 3$. CASE: α is a limit ordinal.
 $\langle 2 \rangle 1$. There exists a set S of ordinals $< \alpha$ such that $|S| = \operatorname{cf} \alpha$ and $\alpha = \operatorname{ssup} S$.
  $\langle 3 \rangle 1$. PICK a set S of ordinals $< \alpha$ such that $|S| = \operatorname{cf} \alpha$ and $\sup S = \alpha$
     PROVE: $\alpha = \operatorname{ssup} S$
  $\langle 3 \rangle 2$. $\forall \beta \in S. \beta < \alpha$
  $\langle 3 \rangle 3$. For any ordinal γ, if $\forall \beta \in S. \beta < \gamma$ then $\alpha \leq \gamma$
 $\langle 2 \rangle 2$. If T is a set of ordinals $< \alpha$ such that $\alpha = \operatorname{ssup} T$, then $\operatorname{cf} \alpha \leq |T|$.
  $\langle 3 \rangle 1$. LET: T be a set of ordinals $< \alpha$ such that $\alpha = \operatorname{ssup} T$
  $\langle 3 \rangle 2$. $\alpha = \sup T$
   $\langle 4 \rangle 1$. For all $\beta \in T$ we have $\beta \leq \alpha$
   $\langle 4 \rangle 2$. LET: μ be any upper bound for T
      PROVE: $\alpha \leq \mu$
   $\langle 4 \rangle 3$. $\alpha \leq \mu + 1$
     PROOF: Since $\forall \beta \in T. \beta < \mu + 1$.
   $\langle 4 \rangle 4$. $\alpha \neq \mu + 1$
     PROOF: Since α is a limit ordinal.
   $\langle 4 \rangle 5$. $\alpha < \mu + 1$
   $\langle 4 \rangle 6$. $\alpha \leq \mu$
  $\langle 3 \rangle 3$. $\operatorname{cf} \alpha \leq |T|$
$\square$

11.5   Inaccessible Cardinals

Definition 11.5.1 (Inaccessible Cardinal)**.** A cardinal number κ is *inaccessible* iff:

- $\kappa > \aleph_0$

- $\forall \lambda < \kappa . 2^\lambda < \kappa$ (cardinal exponentiation)

- κ is regular.

Any inaccessible cardinal is weakly inaccessible.

PROOF:
$\langle 1 \rangle 1.$ LET: $\kappa = \aleph_\lambda$ be weakly inaccessible.
        PROVE: λ is a limit ordinal.
$\langle 1 \rangle 2.$ $\lambda \neq 0$
$\langle 1 \rangle 3.$ ASSUME: for a contradiction $\lambda = \beta + 1$
$\langle 1 \rangle 4.$ $\aleph_\beta < \kappa$
$\langle 1 \rangle 5.$ $2^{\aleph_\beta} < \kappa$
$\langle 1 \rangle 6.$ $\aleph_{\beta+1} < \kappa$
  PROOF: Since $\aleph_{\beta+1} \leq 2^{\aleph_\beta}$.
$\langle 1 \rangle 7.$ Q.E.D.
  PROOF: This contradicts $\langle 1 \rangle 3$.
⬚

Proposition 11.5.2. *If the Generalized Continuum Hypothesis is true, then every weakly inaccessible cardinal is inaccessible.*

PROOF:
$\langle 1 \rangle 1.$ ASSUME: The Generalized Continuum Hypothesis.
$\langle 1 \rangle 2.$ LET: $\kappa = \aleph_\lambda$ be weakly inaccessible.
$\langle 1 \rangle 3.$ $\kappa > \aleph_0$
  PROOF: $\lambda > 0$ because λ is a limit ordinal.
$\langle 1 \rangle 4.$ For all $\mu < \kappa$ we have $2^\mu < \kappa$
  $\langle 2 \rangle 1.$ LET: $\mu < \kappa$
  $\langle 2 \rangle 2.$ LET: $\mu = \aleph_\alpha$
  $\langle 2 \rangle 3.$ $\alpha < \lambda$
  $\langle 2 \rangle 4.$ $\alpha + 1 < \lambda$
    PROOF: λ is a limit ordinal.
  $\langle 2 \rangle 5.$ $2^\mu < \kappa$
    PROOF:

$$
\begin{aligned}
2^\mu &= 2^{\aleph_\alpha} & (\langle 2 \rangle 2) \\
      &= 2^{\beth_\alpha} & (\langle 1 \rangle 1) \\
      &= \beth_{\alpha+1} & \\
      &= \aleph_{\alpha+1} & (\langle 1 \rangle 1) \\
      &< \aleph_\lambda & (\langle 2 \rangle 4) \\
      &= \kappa & (\langle 1 \rangle 2)
\end{aligned}
$$

$\langle 1 \rangle 5$. κ is regular.
    PROOF: $\langle 1 \rangle 2$

□

Lemma 11.5.3. *Let κ be an inaccessible cardinal. For every ordinal $\alpha < \kappa$ we have $\beth_\alpha < \kappa$.*

PROOF:
$\langle 1 \rangle 1$. $\beth_0 < \kappa$
    PROOF: Since $\kappa > \aleph_0$.
$\langle 1 \rangle 2$. For any ordinal α, if $\beth_\alpha < \kappa$ then $\beth_{\alpha+1} < \kappa$.
    PROOF: Since $\beth_{\alpha+1} = 2^{\beth_\alpha} < \kappa$.
$\langle 1 \rangle 3$. For any limit ordinal λ, if $\forall \alpha < \lambda. \beth_\alpha < \kappa$ and $\lambda < \kappa$ then $\beth_\lambda < \kappa$.
    PROOF: By regularity of κ, since $\beth_\lambda$ is the union of $|\lambda|$ cardinals all $< \kappa$.

□

Lemma 11.5.4. *Let κ be an inaccessible cardinal. For all $A \in V_\kappa$ we have $|A| < \kappa$.*

PROOF:
$\langle 1 \rangle 1$. LET: $A \in V_\kappa$
$\langle 1 \rangle 2$. PICK $\alpha < \kappa$ such that $A \in V_\alpha$
$\langle 1 \rangle 3$. $A \subseteq V_\alpha$
$\langle 1 \rangle 4$. $|A| \leq |V_\alpha| \leq \beth_\alpha < \kappa$

□

Theorem Schema 11.5.5. *For every axiom α of ZFC, the following is a theorem:*

    For any inaccessible cardinal κ, we have V_κ is a model of α.

PROOF: For every axiom except the Replacement Axioms, we have Corollary 10.0.11.1.
For an Axiom of Replacement using the property $P[x, y, z_1, \ldots, z_n]$, we reason as follows:
$\langle 1 \rangle 1$. LET: κ be an inaccessible cardinal
      PROVE:
$$\forall a_1, \ldots, a_n, B \in V_\kappa (\forall x \in B. \forall y, y' \in V_\kappa$$
$$(P[x, y, a_1, \ldots, a_n]^{V_\kappa} \wedge P[x, y', a_1, \ldots, a_n]^{V_\kappa} \Rightarrow y = y') \Rightarrow$$
$$\exists C \in V_\kappa \forall y \in V_\kappa (y \in C \Leftrightarrow \exists x \in B. P[x, y, a_1, \ldots, a_n]^{V_\kappa}))$$
$\langle 1 \rangle 2$. LET: $a_1, \ldots, a_n, B \in V_\kappa$
$\langle 1 \rangle 3$. ASSUME: for all $x \in B$, there exists at most one $y \in V_\kappa$ such that $P[x, y, a_1, \ldots, a_n]^{V_\kappa}$.
$\langle 1 \rangle 4$. LET: $F = \{(x, y) \in B \times V_\kappa \mid P[x, y, a_1, \ldots, a_n]^{V_\kappa}\}$
$\langle 1 \rangle 5$. LET: $C = \operatorname{ran} F$
      PROVE: $C \in V_\kappa$
$\langle 1 \rangle 6$. LET: $S = \{\operatorname{rank} F(x) \mid x \in \operatorname{dom} F\}$
$\langle 1 \rangle 7$. $|S| < \kappa$
    PROOF: Since $|S| \leq |\operatorname{dom} F| \leq |B| < \kappa$.

$\langle 1 \rangle 8.$ $\forall \alpha \in S . \alpha < \kappa$
   PROOF: Since $F(x) \in V_\kappa$ for all $x \in \mathrm{dom}\, F$.
$\langle 1 \rangle 9.$ $\sup S < \kappa$
   PROOF: Since κ is regular.
$\langle 1 \rangle 10.$ $\mathrm{rank}\, C \leq \sup S + 1$
$\langle 1 \rangle 11.$ $\mathrm{rank}\, C < \kappa$
$\langle 1 \rangle 12.$ $C \in V_\kappa$
□

Chapter 12

Group Theory

12.1 Groups

Definition 12.1.1 (Group)**.** A *group* G consists of a set G and a function $\cdot : G^2 \to G$ such that:

1. $\cdot$ is associative

2. There exists $e \in G$ such that $\forall x \in G.xe = x$ and $\forall x \in G.\exists y \in G.xy = e$.

Proposition 12.1.2. *The inverse of an element in a group is unique.*

PROOF:
$\langle 1 \rangle 1$. ASSUME: b and b' are inverses of a.
$\langle 1 \rangle 2$. $b = b'$
   PROOF:

$$
\begin{aligned}
b &= be \\
&= bab' \\
&= eb' \\
&= b'
\end{aligned}
$$

$\square$

Definition 12.1.3. We write x^{-1} for the inverse of x.

Proposition 12.1.4. *In any group, if $ab = ac$ then $b = c$.*

PROOF:

$$
\begin{aligned}
b &= eb \\
&= a^{-1}ab \\
&= a^{-1}ac \\
&= ec \\
&= c \qquad\qquad\qquad \square
\end{aligned}
$$

12.2   Abelian Groups

Definition 12.2.1 (Abelian group)**.** An *Abelian group* is a group whose multiplication is commutative.

We may say we are writing an Abelian group *additively*, meaning we write $a + b$ for ab, 0 for e and $-a$ for a^{-1}. In this case we write $a - b$ for ab^{-1}.

Chapter 13

Ring Theory

13.1 Rings

Definition 13.1.1 (Commutative Ring). A *commutative ring* consists of a set R and two binary operations $+, \cdot$ on R such that:

- D is an Abelian group under $+$. Let us write 0 for its identity element.

- $\cdot$ is commutative and associative, and distributes over $+$.

- $\cdot$ has an identity element 1 that is different from 0.

Proposition 13.1.2. *In any commutative ring, $0x = 0$.*

PROOF:
$$(0 + 0)x = 0x$$
$$\therefore 0x + 0x = 0x + 0$$
$$\therefore 0x = 0 \qquad \text{(Proposition 12.1.4)} \square$$

Proposition 13.1.3. *In any commutative ring, $(-a)b = -(ab)$.*

PROOF:
$$ab + (-a)b = (a + (-a))b$$
$$= 0b$$
$$= 0 \qquad \text{(Proposition 13.1.2)} \square$$

13.2 Ordered Rings

Definition 13.2.1 (Ordered Commutative Ring). An *ordered commutative ring* consists of a commutative ring R with a linear order $<$ on R such that:

- for all $x, y, z \in R$, we have $x < y$ if and only if $x + z < y + z$.

- for all $x, y, z \in R$, if $0 < z$ then we have $x < y$ if and only if $xz < yz$.

Proposition 13.2.2. *In any ordered commutative ring, $0 < 1$.*

PROOF: If $1 < 0$ then we have $0 < -1$ and so $0 < (-1)(-1) = 1$, which is a contradiction. □

Proposition 13.2.3. *The ordering on an ordered commutative ring is dense; that is, if $x < y$ then there exists z such that $x < z < y$.*

PROOF: Take $z = (x + y)/2$. □

13.3   Integral Domains

Definition 13.3.1 (Integral Domain). An *integral domain* is a commutative ring such that, for all $a, b \in D$, if $ab = 0$ then $a = 0$ or $b = 0$.

Proposition 13.3.2. *In any integral domain, if $ab = ac$ and $a \neq 0$ then $b = c$.*

PROOF: We have $a(b - c) = 0$ and $a \neq 0$ so $b - c = 0$ hence $b = c$. □

Definition 13.3.3 (Ordered Integral Domain). An *ordered integral domain* is an ordered commutative ring that is an integral domain.

Chapter 14

Field Theory

14.1 Fields

Definition 14.1.1 (Field)**.** A *field* F is a commutative ring such that $0 \neq 1$ and, for all $x \in F$, if $x \neq 0$ then there exists $y \in F$ such that $xy = 1$.

Proposition 14.1.2. *Every field is an integral domain.*

PROOF: If $ab = 0$ and $a \neq 0$ then $b = a^{-1}ab = 0$. □

Proposition 14.1.3. *In any field F, we have $F - \{0\}$ is an Abelian group under multiplication.*

PROOF: Immediate from the definition. □

Definition 14.1.4 (Field of Fractions)**.** Let D be an integral domain. The *field of fractions* of D is the quotient set $F = (D \times (D - \{0\}))/\sim$ where

$$(a, b) \sim (c, d) \Leftrightarrow ad = bc$$

under

$$[(a, b)] + [(c, d)] = [(ad + bc, bd)]$$
$$[(a, b)][(c, d)] = [(ac, bd)]$$

We prove this is a field.

PROOF:
⟨1⟩1. $\sim$ is an equivalence relation on $D \times (D - \{0\})$.
  PROOF:
    ⟨2⟩1. $\sim$ is reflexive.
      PROOF: We always have $ab = ba$.
    ⟨2⟩2. $\sim$ is symmetric.
      PROOF: If $ad = bc$ then $cb = da$.

$\langle 2 \rangle 3.$ $\sim$ is transitive.

   $\langle 3 \rangle 1.$ ASSUME: $(a, b) \sim (c, d) \sim (e, f)$

   $\langle 3 \rangle 2.$ $ad = bc$ and $cf = de$

   $\langle 3 \rangle 3.$ $adf = bde$

     PROOF: $adf = bcf = bde$

   $\langle 3 \rangle 4.$ $af = be$

     PROOF: Proposition 13.3.2.

$\square$

$\langle 1 \rangle 2.$ Addition is well-defined.

  PROOF:

  $\langle 2 \rangle 1.$ If $b \neq 0$ and $d \neq 0$ then $bd \neq 0$.

   PROOF: Since D is an integral domain.

  $\langle 2 \rangle 2.$ If $ab' = a'b$ and $cd' = c'd$ then $(ad + bc)b'd' = (a'd' + b'c')bd$.

   PROOF:

$$(ad + bc)b'd' = ab'dd' + bb'cd'$$
$$= a'bdd' + bb'c'd$$
$$= (a'd' + b'c')bd$$

$\square$

$\langle 1 \rangle 3.$ Multiplication is well-defined.

  PROOF:

  $\langle 2 \rangle 1.$ If $b \neq 0$ and $d \neq 0$ then $bd \neq 0$.

   PROOF: Since D is an integral domain.

  $\langle 2 \rangle 2.$ If $[(a, b)] = [(a', b')]$ and $[(c, d)] = [(c', d')]$ then $[(ac, bd)] = [(a'c', b'd')]$.

   PROOF: If $ab' = a'b$ and $cd' = c'd$ then $acb'd' = a'c'bd$.

$\square$

$\langle 1 \rangle 4.$ Addition is commutative.

  PROOF: $[(a, b)] + [(c, d)] = [(ad + bc, bd)] = [(cb + da, db)] = [(c, d)] + [(a, b)]$ $\square$

$\langle 1 \rangle 5.$ Addition is associative.

  PROOF:

$$[(a, b)] + ([(c, d)] + [(e, f)]) = [(a, b)] + [(cf + de, df)]$$
$$= [(adf + bcf + bde, bdf)]$$
$$= [(ad + bc, bd)] + [(e, f)]$$
$$= ([(a, b)] + [(c, d)]) + [(e, f)] \qquad \square$$

$\langle 1 \rangle 6.$ For any $x \in F$ we have $x + [(0, 1)] = x$

  PROOF: $[(a, b)] + [(0, 1)] = [(a \cdot 1 + b \cdot 0, b \cdot 1)] = [(a, b)]$ $\square$

$\langle 1 \rangle 7.$ For any $x \in F$, there exists $y \in F$ such that $x + y = [(0, 1)]$.

  PROOF: $[(a, b)] + [(-a, b)] = [(ab - ab, b^2)] = [(0, b^2)] = [(0, 1)]$ $\square$

$\langle 1 \rangle 8.$ Multiplication is commutative.

  PROOF: $[(a, b)][(c, d)] = [(c, d)][(a, b)] = [(ac, bd)]$. $\square$

$\langle 1 \rangle 9.$ Multiplication is assocative.

  PROOF: $[(a, b)]([(c, d)][(e, f)]) = ([(a, b)][(c, d)])[(e, f)] = [(ace, bdf)]$. $\square$

$\langle 1 \rangle 10.$ For any $x \in F$ we have $x[(1, 1)] = x$

  PROOF: $[(a, b)][(1, 1)] = [(a, b)]$ $\square$

$\langle 1 \rangle 11.$ For any non-zero $x \in F$, there exists $y \in F$ such that $xy = [(1, 1)]$.

PROOF:

⟨2⟩1. LET: $[(a,b)] \in \mathbb{Q}$

⟨2⟩2. ASSUME: $[(a,b)] \neq [(0,1)]$

⟨2⟩3. $a \neq 0$

⟨2⟩4. $[(a,b)][(b,a)] = [(1,1)]$

□

□

Definition 14.1.5. For any field F, let $N(F)$ be the intersection of all the subsets $S \subseteq F$ such that $1 \in S$ and $\forall x \in S.x + 1 \in S$.

Definition 14.1.6 (Characteristic Zero)**.** A field F has *characteristic 0* iff $0 \notin N(F)$.

Proposition 14.1.7. *In a field F with characteristic 0, the function $n : \mathbb{N} \to N(F)$ defined by*

$$n(0) = 1$$
$$n(x + 1) = n(x) + 1$$

is a bijection.

PROOF:

⟨1⟩1. n is injective.

  ⟨2⟩1. ASSUME: for a contradiction $n(i) = n(j)$ with $i \neq j$

  ⟨2⟩2. ASSUME: w.l.o.g. $i < j$

  ⟨2⟩3. $n(j - i) = 0$

  ⟨2⟩4. Q.E.D.

    PROOF: This contradicts the fact that F has characteristic 0.

⟨1⟩2. n is surjective.

  PROOF: Since ran n is a subset of F that includes 1 and is closed under $+1$.

□

Definition 14.1.8. In any field F, let

$$I(F) = N(F) \cup \{0\} \cup \{-x \mid x \in N(F)\}$$

Definition 14.1.9. In any field F, let

$$Q(F) = \{x/y \mid x, y \in I(F), y \neq 0\}$$

Proposition 14.1.10. $Q(F)$ *is the smallest subfield of F.*

PROOF: $Q(F)$ is closed under $+$ and $\cdot$, and any subset of F closed under $+$ and $\cdot$ that contains 0 and 1 must include $Q(F)$. □

Theorem 14.1.11. *Let F and G be fields of characteristic 0. Then there exists a unique field isomorphism between $Q(F)$ and $Q(G)$.*

PROOF:

⟨1⟩1. LET: $\phi : N(F) \to N(G)$ be the unique function such that $\phi(1) = 1$ and
        $\forall x \in N(F).\phi(x + 1) = \phi(x) + 1$.

⟨1⟩2. ϕ is a bijection.
    PROOF: Similar to Proposition 14.1.7.

⟨1⟩3. $\forall x, y \in N(F).\phi(x + y) = \phi(x) + \phi(y)$
    PROOF: Induction on y.

⟨1⟩4. $\forall x, y \in N(F).\phi(xy) = \phi(x)\phi(y)$
    PROOF: Induction on y.

⟨1⟩5. Extend ϕ to a bijection $I(F) \cong I(G)$ such that $\forall x, y \in I(F).\phi(x + y) =$
        $\phi(x) + \phi(y)$ and $\forall x, y \in I(F).\phi(xy) = \phi(x)\phi(y)$

  ⟨2⟩1. Define $\phi(0) = 0$ and $\phi(-x) = -\phi(x)$ for $x \in N(F)$

    ⟨3⟩1. $0 \notin N(F)$

    ⟨3⟩2. For all $x \in N(F)$ we have $-x \notin N(F)$
      PROOF: Then we would have $x + -x = 0 \in N(F)$.

    ⟨3⟩3. For all $x \in N(F)$ we have $-x \neq 0$

  ⟨2⟩2. For all $x, y \in I(F)$ we have $\phi(x + y) = \phi(x) + \phi(y)$
      PROOF: Case analysis on x and y.

  ⟨2⟩3. For all $x, y \in I(F)$ we have $\phi(xy) = \phi(x)\phi(y)$
      PROOF: Case analysis on x and y.

⟨1⟩6. Extend ϕ to a bijection $Q(F) \cong Q(G)$ such that $\forall x, y \in Q(F).\phi(x + y) =$
        $\phi(x) + \phi(y)$ and $\forall x, y \in Q(F).\phi(xy) = \phi(x)\phi(y)$

  ⟨2⟩1. Define $\phi(x/y) = \phi(x)/\phi(y)$

⟨1⟩7. ϕ is unique.

  ⟨2⟩1. LET: θ satisfy the theorem.

  ⟨2⟩2. For all $x \in N(F)$ we have $\theta(x) = \phi(x)$

  ⟨2⟩3. For all $x \in I(F)$ we have $\theta(x) = \phi(x)$

  ⟨2⟩4. For all $x \in Q(F)$ we have $\theta(x) = \phi(x)$

□

14.2   Ordered Fields

Definition 14.2.1 (Ordered Field). An *ordered field* is an ordered commutative ring that is a field.

Proposition 14.2.2. *Every ordered field F has characteristic 0.*

PROOF: We have $0 < n$ for all $n \in N(F)$. □

Proposition 14.2.3. *Let F be a field of characteristic 0. Then there exists a unique relation $<$ on $Q(F)$ that makes $Q(F)$ into an ordered field.*

PROOF: Easy. □

Corollary 14.2.3.1. *Let F and G be ordered fields. Let ϕ be the unique field isomorphism between $Q(F)$ and $Q(G)$. Then ϕ is an ordered field isomorphism.*

Definition 14.2.4 (Archimedean). An ordered field F is *Archimedean* iff

$$\forall x \in F.\exists n \in N(F).n > x \ .$$

Proposition 14.2.5. *Let F be an Archimedean ordered field. Let $x, y \in F$ with $x > 0$. Then there exists $n \in N(F)$ such that $nx > y$.*

PROOF: Pick $n > y/x$. □

Proposition 14.2.6. *Let F be an Archimedean ordered field. For all $x, y \in F$, if $x < y$, then there exists $r \in Q(F)$ such that $x < r < y$.*

PROOF:
⟨1⟩1. CASE: $x > 0$
  ⟨2⟩1. PICK $n \in N(F)$ such that $n(y - x) > 1$
    PROOF: Proposition 14.2.5.
  ⟨2⟩2. $ny > 1 + nx$
  ⟨2⟩3. LET: m be the least element of $N(F)$ such that $m > nx$.
  ⟨2⟩4. $m - 1 \le nx$
  ⟨2⟩5. $nx < m < ny$
  ⟨2⟩6. $x < m/n < y$
⟨1⟩2. CASE: $x \le 0$
  ⟨2⟩1. PICK $k \in N(F)$ such that $k > -x$
  ⟨2⟩2. $0 < x + k < y + k$
  ⟨2⟩3. PICK $r \in Q(F)$ such that $x + k < r < y + k$
    PROOF: ⟨1⟩1
  ⟨2⟩4. $x < r - k < y$

Definition 14.2.7 (Complete)**.** An ordered field F is *complete* iff every nonempty subset of F bounded above has a least upper bound.

Proposition 14.2.8. *Every complete ordered field is Archimedean.*

PROOF:
⟨1⟩1. LET: F be a complete ordered field.
⟨1⟩2. LET: $x \in F$
⟨1⟩3. ASSUME: for a contradiction there is no member of $N(F)$ greater than x.
⟨1⟩4. x is an upper bound for $N(F)$.
⟨1⟩5. LET: $y = \sup N(F)$
⟨1⟩6. PICK $n \in N(F)$ such that $y - 1 < n$
⟨1⟩7. $y < n + 1$
⟨1⟩8. Q.E.D.
  PROOF: This is a contradiction.
□

Proposition 14.2.9. *Let F be a complete ordered field and $a \in F$ be nonnegative. Then there exists $b \in F$ such that $b^2 = a$.*

PROOF:
⟨1⟩1. LET: $B = \{x \in F \mid 0 \le x \le 1 + a\}$
⟨1⟩2. LET: $\phi : B \to B$ be the function

$$\phi(x) = x + \frac{1}{2(1 + a)}(a - x^2) \ .$$

$\langle 1\rangle 3$. ϕ is strictly monotone.

   $\langle 2\rangle 1$. LET: $0 \le x < y \le 1 + a$

   $\langle 2\rangle 2$. $1 - \frac{x+y}{2(1+a)} > 0$

   $\langle 2\rangle 3$. $\phi(y) - \phi(x) = (y - x)(1 - \frac{x+y}{2(1+a)}) > 0$

   $\langle 2\rangle 4$. $\phi(x) < \phi(y)$

$\langle 1\rangle 4$. PICK $b \in B$ such that $\phi(b) = b$.

   PROOF: Knaster Fixed-Point Theorem.

$\langle 1\rangle 5$. $b^2 = a$

$\square$

Theorem 14.2.10 (Uniqueness of the Complete Ordered Field). *If F and G are complete ordered fields, then there exists a unique bijection $\phi : F \cong G$ such that, for all $x, y \in F$,*

$$\phi(x + y) = \phi(x) + \phi(y)$$
$$\phi(xy) = \phi(x)\phi(y)$$

This bijection also satisfies: for all $x, y \in F$,

$$x < y \Leftrightarrow \phi(x) < \phi(y) \ .$$

PROOF:

$\langle 1\rangle 1$. PICK a bijection $\phi : Q(F) \cong Q(G)$ such that, for all $x, y \in Q(F)$,
$$\phi(x + y) = \phi(x) + \phi(y)$$
$$\phi(xy) = \phi(x)\phi(y)$$
$$x < y \Leftrightarrow \phi(x) < \phi(y)$$

   PROOF: Corollary 14.2.3.1.

$\langle 1\rangle 2$. $Q(F)$ intersects every interval in F.

   PROOF: Proposition 14.2.6.

$\langle 1\rangle 3$. $Q(G)$ intersects every interval in G.

   PROOF: Proposition 14.2.6.

$\langle 1\rangle 4$. PICK an order isomorphism $\psi : F \cong G$ that extends ϕ.

   PROOF: Theorem 5.1.21.

$\langle 1\rangle 5$. $\forall x, y \in F.\psi(x + y) = \psi(x) + \psi(y)$

   $\langle 2\rangle 1$. LET: $x, y \in F$

   $\langle 2\rangle 2$. $\psi(x) + \psi(y) \not< \psi(x + y)$

      $\langle 3\rangle 1$. ASSUME: for a contradiction $\psi(x) + \psi(y) < \psi(x + y)$

      $\langle 3\rangle 2$. PICK $r' \in Q(G)$ such that $\psi(x) < r' < \psi(x + y) - \psi(y)$

      $\langle 3\rangle 3$. PICK $s' \in Q(G)$ such that $\psi(y) < s' < \psi(x + y) - r'$

      $\langle 3\rangle 4$. $r' + s' < \psi(x + y)$

      $\langle 3\rangle 5$. PICK $r, s \in Q(F)$ such that $\phi(r) = r'$ and $\phi(s) = s'$

      $\langle 3\rangle 6$. $\phi(r + s) = r' + s'$

      $\langle 3\rangle 7$. $\psi(x) < \psi(r)$

      $\langle 3\rangle 8$. $\psi(y) < \psi(s)$

      $\langle 3\rangle 9$. $\psi(x + y) > \psi(r + s)$

      $\langle 3\rangle 10$. $x < r$

⟨3⟩11. $y < s$
⟨3⟩12. $x + y > r + s$
⟨3⟩13. Q.E.D.
PROOF: This is a contradiction.
⟨2⟩3. $\psi(x + y) \not< \psi(x) + \psi(y)$
PROOF: Similar.
⟨1⟩6. $\forall x, y \in F.\psi(xy) = \psi(x)\psi(y)$
⟨2⟩1. LET: $x, y \in F$
⟨2⟩2. CASE: x and y are positive.
⟨3⟩1. $\psi(x)\psi(y) \not< \psi(xy)$
⟨4⟩1. ASSUME: for a contradiction $\psi(x)\psi(y) < \psi(xy)$
⟨4⟩2. PICK $r' \in Q(G)$ such that $\psi(x) < r' < \psi(xy)/\psi(y)$
⟨4⟩3. PICK $s' \in Q(G)$ such that $\psi(y) < s' < \psi(xy)/r'$
⟨4⟩4. $r's' < \psi(xy)$
⟨4⟩5. PICK $r, s \in Q(F)$ such that $\phi(r) = r'$ and $\phi(s) = s'$
⟨4⟩6. $\phi(rs) = r's'$
⟨4⟩7. $x < r$, $y < s$ and $rs < xy$
⟨4⟩8. Q.E.D.
PROOF: This is a contradiction.
⟨3⟩2. $\psi(xy) \not< \psi(x)\psi(y)$
PROOF: Similar.
⟨2⟩3. CASE: x and y are not both positive.
PROOF: Follows from ⟨2⟩2 since $\psi(-x) = -\psi(x)$ by ⟨1⟩5.
⟨1⟩7. For any field isomorphism $\theta : F \cong G$, we have $\theta = \psi$.
⟨2⟩1. $\theta \upharpoonright Q(F) = \phi$
PROOF: Theorem 14.1.11.
⟨2⟩2. θ is strictly monotone.
⟨3⟩1. LET: $x, y \in F$ with $x < y$
⟨3⟩2. $y - x > 0$
⟨3⟩3. PICK $z \in F$ such that $z^2 = y - x$
⟨3⟩4. $\theta(z)^2 = \theta(y) - \theta(x)$
⟨3⟩5. $\theta(y) - \theta(x) > 0$
⟨3⟩6. $\theta(x) < \theta(y)$
⟨2⟩3. $\theta = \psi$
PROOF: By the uniqueness of ψ.
☐

Chapter 15

Number Systems

15.1 The Integers

Definition 15.1.1. The set of *integers* $\mathbb{Z}$ is the quotient set $\mathbb{N}^2/\sim$, where $(m, n) \sim (p, q)$ iff $m + q = n + p$.

We prove $\sim$ is an equivalence relation on $\mathbb{N}^2$.

PROOF:
$\langle 1 \rangle 1.$ $\sim$ is reflexive.
    PROOF: For all $m, n \in \mathbb{N}$ we have $m + n = n + m$.
$\langle 1 \rangle 2.$ $\sim$ is symmetric.
    PROOF: If $m + q = n + p$ then $p + n = q + m$.
$\langle 1 \rangle 3.$ $\sim$ is transitive.
    $\langle 2 \rangle 1.$ ASSUME: $(m, n) \sim (p, q) \sim (r, s)$
    $\langle 2 \rangle 2.$ $m + q = n + p$ and $p + s = q + r$
    $\langle 2 \rangle 3.$ $m + q + s = n + q + r$
    $\langle 2 \rangle 4.$ $m + s = n + r$
        PROOF: By cancellation.
☐

Definition 15.1.2 (Addition). Define *addition* $+$ on $\mathbb{Z}$ by $[(m, n)] + [(p, q)] = [(m + p, n + q)]$.

We prove this is well-defined.

PROOF: If $m + n' = n + m'$ and $p + q' = q + p'$ then $m + p + n' + q' = n + q + m' + p'$.
☐

Proposition 15.1.3. *Addition on $\mathbb{Z}$ is commutative.*

PROOF: $[(m, n)] + [(p, q)] = [(m + p, n + q)] = [(p + m, q + n)] = [(p, q)] + [(m, n)]$.
☐

Proposition 15.1.4. *Addition on $\mathbb{Z}$ is associative.*

PROOF: $[(m,n)] + ([(p,q)] + [(r,s)]) = [(m+p+r, n+q+s)] = ([(m,n)] + [(p,q)]) + [(r,s)]$. □

Proposition 15.1.5. *Given natural numbers m and n, we have $[(m,0)] = [(n,0)]$ iff $m = n$.*

PROOF: Immediate from definitions. □

Definition 15.1.6. We identify any natural number n with the integer $[(n,0)]$.

Proposition 15.1.7. *Addition on integers agrees with addition on natural numbers.*

PROOF: Since $[(m,0)] + [(n,0)] = [(m+n,0)]$. □

Proposition 15.1.8. *For all $a \in \mathbb{Z}$ we have $a + 0 = a$.*

PROOF: $[(m,n)] + [(0,0)] = [(m+0, n+0)] = [(m,n)]$. □

Proposition 15.1.9. *For all $a \in \mathbb{Z}$, there exists $b \in \mathbb{Z}$ such that $a + b = 0$.*

PROOF: $[(m,n)] + [(n,m)] = [(m+n, m+n)] = [(0,0)]$ □

Proposition 15.1.10. *The integers form an Abelian group under addition.*

PROOF: Proposition 15.1.3, 15.1.4, 15.1.8, 15.1.9. □

Definition 15.1.11. Define multiplication $\cdot$ on $\mathbb{Z}$ by: $[(m,n)][(p,q)] = [(mp + nq, mq + np)]$.
    We prove this is well defined.

PROOF:
$\langle 1 \rangle 1.$ ASSUME: $m + n' = n + m'$ and $p + q' = q + p'$
        PROVE: $mp + nq + m'q' + n'p' = mq + np + m'p' + n'q'$
$\langle 1 \rangle 2.$ $mp + n'p = np + m'p$
$\langle 1 \rangle 3.$ $nq + m'q = mq + n'q$
$\langle 1 \rangle 4.$ $m'p + m'q' = m'q + m'p'$
$\langle 1 \rangle 5.$ $n'q + n'p' = n'p + n'q'$
$\langle 1 \rangle 6.$ $mp + n'p + nq + m'q + m'p + m'q' + n'q + n'p' = np + m'p + mq + n'q + m'q + m'p' + n'p + n'q'$
$\langle 1 \rangle 7.$ $mp + nq + m'q' + n'p' = mq + np + m'p' + n'q'$
    PROOF: By cancellation.
□

Proposition 15.1.12. *Multiplication on integers agrees with multiplication on natural numbers.*

PROOF: Since $[(m,0)][(n,0)] = [(mn + 0, m0 + n0)] = [(mn, 0)]$. □

Proposition 15.1.13. *Multiplication on $\mathbb{Z}$ is commutative.*

PROOF: $[(m,n)][(p,q)] = [(mp + nq, mq + np)] = [(pm + qn, pn + qm)] = [(p,q)][(m,n)]$. □

Proposition 15.1.14. *Multiplication on $\mathbb{Z}$ is associative.*

Proof:
$$
\begin{aligned}
[(m,n)]([(p,q)][(r,s)]) &= [(m,n)][(pr+qs, ps+qr)] \\
&= [(mpr+mqs+nps+nqr, mps+mqr+npr+nqs)] \\
&= [(mp+nq, mq+np)][(r,s)] \\
&= ([(m,n)][(p,q)])[(r,s)] \qquad\qquad \square
\end{aligned}
$$

Proposition 15.1.15. *Multiplication distributes over addition.*

Proof:
$$
\begin{aligned}
[(m,n)]([(p,q)] + [(r,s)]) &= [(m,n)][(p+r, q+s)] \\
&= [(mp+mr+nq+ns, np+nr+mq+ms)] \\
[(m,n)][(p,q)] + [(m,n)][(r,s)] &= [(mp+nq, mq+np)] + [(mr+ns, ms+nr)] \\
&= [(mp+nq+mr+ns, mq+np+ms+nr)] \qquad \square
\end{aligned}
$$

Proposition 15.1.16. *For any integer a we have $a1 = a$.*

Proof: Since $[(m,n)][(1,0)] = [(m1+n0, m0+n1)] = [(m,n)]$. $\square$

Proposition 15.1.17. *For any integeres a and b, if $ab = 0$ then $a = 0$ or $b = 0$.*

Proof:
$\langle 1 \rangle$1. Assume: $[(m,n)][(p,q)] = [(0,0)]$
$\langle 1 \rangle$2. $mp + nq = mq + np$
$\langle 1 \rangle$3. Assume: $[(m,n)] \neq [(0,0)]$
$\langle 1 \rangle$4. $m \neq n$
        Prove: $p = q$
$\langle 1 \rangle$5. Case: $m < n$
   $\langle 2 \rangle$1. $p \not< q$
      Proof: If $p < q$ then $mq + np < mp + nq$ by Proposition 8.4.6.
   $\langle 2 \rangle$2. $q \not< p$
      Proof: If $q < p$ then $mp + nq < mq + np$ by Proposition 8.4.6.
   $\langle 2 \rangle$3. $p = q$
      Proof: By trichotomy.
$\langle 1 \rangle$6. Case: $n < m$
   Proof: Similar.
$\square$

Proposition 15.1.18. *The integers $\mathbb{Z}$ form an integral domain.*

Proof: Propositions 15.1.13, 15.1.14, 15.1.15, 15.1.16, 15.1.17, 15.1.10. $\square$

Definition 15.1.19. Define $<$ on $\mathbb{Z}$ by $[(m,n)] < [(p,q)]$ if and only if $m + q < n + p$.
    We prove this is well-defined.

Proof:

$\langle 1 \rangle 1$. ASSUME: $m + n' = n + m'$ and $p + q' = q + p'$.
        PROVE: $m + q < n + p$ if and only if $m' + q' < n' + p'$
$\langle 1 \rangle 2$. $m + q < n + p$ if and only if $m' + q' < n' + p'$
    PROOF:

$$
\begin{aligned}
m + q < n + p &\Leftrightarrow m + n' + q < n + n' + p && \text{(Corollary 6.5.7.1)} \\
&\Leftrightarrow m' + n + q < n + n' + p && \\
&\Leftrightarrow m' + q < n' + p && \text{(Corollary 6.5.7.1)} \\
&\Leftrightarrow m' + q + p' < n' + p + p' && \text{(Corollary 6.5.7.1)} \\
&\Leftrightarrow m' + q' + p < n' + p + p' && \\
&\Leftrightarrow m' + q' < n' + p' && \text{(Corollary 6.5.7.1)}\ \square
\end{aligned}
$$

Proposition 15.1.20. *The ordering on the integers agrees with the ordering on the natural numbers.*

PROOF: We have $[(m, 0)] < [(n, 0)]$ iff $m < n$. $\square$

Proposition 15.1.21. $<$ *is a linear order on* $\mathbb{Z}$.

PROOF:
$\langle 1 \rangle 1$. $<$ is irreflexive.
    PROOF: We never have $m + n < m + n$.
$\langle 1 \rangle 2$. $<$ is transitive.
    $\langle 2 \rangle 1$. ASSUME: $[(m, n)] < [(p, q)] < [(r, s)]$
    $\langle 2 \rangle 2$. $m + q < n + p$ and $p + s < q + r$
    $\langle 2 \rangle 3$. $m + q + s < n + q + r$
        PROOF: $m + q + s < n + p + s < n + q + r$
    $\langle 2 \rangle 4$. $m + s < n + r$
        PROOF: Corollary 6.5.7.1.
$\langle 1 \rangle 3$. $<$ is total.
    PROOF: Given natural numbers m, n, p and q, either $m + q < n + p$, or $m + q = n + p$, or $n + p < m + q$.
$\square$

Definition 15.1.22 (Positive). An integer a is *positive* iff $a > 0$.

Theorem 15.1.23. *For any integers a, b and c, we have $a < b$ if and only if $a + c < b + c$.*

PROOF:
$\langle 1 \rangle 1$. If $a < b$ then $a + c < b + c$.
    $\langle 2 \rangle 1$. LET: $a = [(m, n)]$, $b = [(p, q)]$ and $c = [(r, s)]$.
    $\langle 2 \rangle 2$. ASSUME: $a < b$
    $\langle 2 \rangle 3$. $m + q < n + p$
    $\langle 2 \rangle 4$. $m + r + q + s < n + r + p + s$
    $\langle 2 \rangle 5$. $[(m + r, n + s)] < [(p + r, q + s)]$
    $\langle 2 \rangle 6$. $a + c < b + c$

$\langle 1 \rangle 2$. If $a + c < b + c$ then $a < b$.
  PROOF: From $\langle 1 \rangle 1$ and Proposition 5.2.6.
□

Proposition 15.1.24. *Let a, b and c be integers. If $0 < c$, then $a < b$ if and only if $ac < bc$.*

PROOF:
$\langle 1 \rangle 1$. LET: $c = [(r, s)]$
$\langle 1 \rangle 2$. ASSUME: $0 < c$
$\langle 1 \rangle 3$. $s < r$
$\langle 1 \rangle 4$. For all integers a and b, if $a < b$ then $ac < bc$
  $\langle 2 \rangle 1$. LET: $a = [(m, n)]$, $b = [(p, q)]$.
  $\langle 2 \rangle 2$. ASSUME: $a < b$
  $\langle 2 \rangle 3$. $m + q < n + p$
  $\langle 2 \rangle 4$. $(m + q)r + (p + n)s < (m + q)s + (p + n)r$
    PROOF: Proposition 8.4.6, $\langle 1 \rangle 3$, $\langle 2 \rangle 3$.
  $\langle 2 \rangle 5$. $mr + ns + ps + qr < ms + nr + pr + qs$
  $\langle 2 \rangle 6$. $[(mr + ns, ms + nr)] < [(pr + qs, ps + qr)]$
  $\langle 2 \rangle 7$. $ac < bc$
$\langle 1 \rangle 5$. For all integers a and b, if $ac < bc$ then $a < b$
  PROOF: From $\langle 1 \rangle 4$ and Proposition 5.2.6.
□

Proposition 15.1.25. *Let a be a positive integer. For any integer b, there exists $k \in \mathbb{N}$ such that $b < ak$.*

PROOF:
$\langle 1 \rangle 1$. CASE: $b \leq 0$
  PROOF: Take $k = 1$.
$\langle 1 \rangle 2$. CASE: $b > 0$
  PROOF: Take $k = b + 1$.
□

15.2 The Rationals

Definition 15.2.1 (Rational Numbers)**.** The set $\mathbb{Q}$ of *rational numbers* is the field of fractions over the integers.

Proposition 15.2.2. *For any integers a and b, we have $[(a, 1)] = [(b, 1)]$ iff $a = b$.*

PROOF: Immediate from definitions. □

  Henceforth we identify any integer a with the rational number $[(a, 1)]$.

Proposition 15.2.3. *Addition on the rationals agrees with addition on the integers.*

PROOF: $[(a, 1)] + [(b, 1)] = [(a \cdot 1 + b \cdot 1, 1 \cdot 1)] = [(a + b, 1)]$. □

Proposition 15.2.4. *Multiplication on the rationals agrees with multiplication on the integers.*

PROOF: $[(a, 1)][(b, 1)] = [(ab, 1)]$ □

Definition 15.2.5. Define the ordering $<$ on the rationals by: if b and d are positive, then $[(a, b)] < [(c, d)]$ iff $ad < bc$.
    We prove this is well-defined.

PROOF:
$\langle 1 \rangle 1$. For any rational q, there exist integers a, b with b positive such that
        $q = [(a, b)]$.
    PROOF: Since $[(a, b)] = [(-a, -b)]$, and if $b \neq 0$ then one of b and $-b$ is
    positive.
$\langle 1 \rangle 2$. If b, b', d and d' are positive, $[(a, b)] = [(a', b')]$, and $[(c, d)] = [(c', d')]$,
        then $ad < bc$ iff $a'd' < b'c'$.
    PROOF:
    $\langle 2 \rangle 1$. If $ad < bc$ then $a'd' < b'c'$.
        $\langle 3 \rangle 1$. ASSUME: $ad < bc$
        $\langle 3 \rangle 2$. $ab'd < bb'c$
        $\langle 3 \rangle 3$. $a'bd < bb'c$
        $\langle 3 \rangle 4$. $a'd < b'c$
        $\langle 3 \rangle 5$. $a'dd' < b'cd'$
        $\langle 3 \rangle 6$. $a'dd' < b'c'd$
        $\langle 3 \rangle 7$. $a'd' < b'c'$
    $\langle 2 \rangle 2$. If $a'd' < b'c'$ then $ad < bc$.
        PROOF: Similar.
□

Proposition 15.2.6. *The ordering on the rationals agrees with the ordering on the integers.*

PROOF: We have $[(a, 1)] < [(b, 1)]$ if and only if $a < b$. □

Proposition 15.2.7. *The relation $<$ is a linear ordering on $\mathbb{Q}$.*

PROOF:
$\langle 1 \rangle 1$. $<$ is irreflexive.
    PROOF: We never have $ab < ab$.
$\langle 1 \rangle 2$. $<$ is transitive.
    $\langle 2 \rangle 1$. ASSUME: $[(a, b)] < [(c, d)] < [(e, f)]$ where b, d and f are positive.
    $\langle 2 \rangle 2$. $ad < bc$ and $cf < de$
    $\langle 2 \rangle 3$. $adf < bde$
        PROOF: $adf < bcf < bde$
    $\langle 2 \rangle 4$. $af < be$
$\langle 1 \rangle 3$. $<$ is total.
    PROOF: For any integers a, b, c, d, we have $ad < bc$ or $ad = bc$ or $bc < ad$.

□

Proposition 15.2.8. *For any rationals r, s and t, we have $r < s$ if and only if $r + t < s + t$.*

PROOF:
⟨1⟩1. LET: a, b, c, d, e, f be integers with b, d and f positive.
⟨1⟩2. $[(a,b)] + [(e,f)] < [(c,d)] + [(e,f)]$ if and only if $[(a,b)] < [(c,d)]$.
  PROOF:
$$\begin{aligned}
[(a,b)] + [(e,f)] < [(c,d)] + [(e,f)] &\Leftrightarrow [(af + be, bf)] < [(cf + de, df)] \\
&\Leftrightarrow (af + be)df < (cf + de)bf \\
&\Leftrightarrow afdf + bedf < cfbf + debf \\
&\Leftrightarrow afdf < cfbf \\
&\Leftrightarrow ad < bc \\
&\Leftrightarrow [(a,b)] < [(c,d)]
\end{aligned}$$

□

Corollary 15.2.8.1. *For any rational r, we have $r < 0$ if and only if $0 < -r$.*

Definition 15.2.9 (Absolute Value)**.** For any rational r, the *absolute value* of r is defined by
$$|r| := \begin{cases} -r & \text{if } 0 < -r \\ r & \text{otherwise} \end{cases}$$

Proposition 15.2.10. *For any rationals r, s and t, if t is positive then $r < s$ iff $rt < st$.*

PROOF:
⟨1⟩1. LET: $r = [(a,b)]$, $s = [(c,d)]$ and $t = [(e,f)]$ where b, d and f are positive.
⟨1⟩2. ASSUME: $0 < t$
⟨1⟩3. $e > 0$
⟨1⟩4. $rt < st$ iff $r < s$
  PROOF:
$$\begin{aligned}
rt < st &\Leftrightarrow [(ae, bf)] < [(ce, df)] \\
&\Leftrightarrow aedf < cebf \\
&\Leftrightarrow ad < bc \\
&\Leftrightarrow r < s
\end{aligned}$$

□

Corollary 15.2.10.1. *The rationals form an ordered field.*

Proposition 15.2.11. *Let p be a positive rational. For any rational number r, there exists $k \in \mathbb{N}$ such that $r < pk$.*

PROOF:
⟨1⟩1. LET: $p = a/b$ and $r = c/d$ where a, b and d are positive.

$\langle 1 \rangle 2$. PICK $k \in \mathbb{N}$ such that $bc < adk$
    PROOF: Proposition 15.1.25.
$\langle 1 \rangle 3$. $r < pk$
▯

Proposition 15.2.12. $\mathbb{Q} \approx \mathbb{N}$

PROOF: Arrange the rationals in order $0/1$, $1/1$, $1/2$, $0/2$, $-1/2$, $-1/1$, $-2/1$, $-2/2$, $-2/3$, $-1/3$, $0/3$, $1/3$, $2/3$, etc. then remove all duplicates. ▯

15.3   The Real Numbers

Definition 15.3.1 (Cauchy Sequence)**.** A *Cauchy sequence* is a sequence (q_n) of rationals such that, for every positive rational ϵ, there exists $k \in \mathbb{N}$ such that $\forall m, n > k.|q_m - q_n| < \epsilon$.

Definition 15.3.2 (Dedekind Cut)**.** A *Dedekind cut* is a set $x \subseteq \mathbb{Q}$ such that:

1. $\emptyset \neq x \neq \mathbb{Q}$

2. x is closed downwards.

3. x has no greatest member.

    The set $\mathbb{R}$ of *real numbers* is the set of Dedekind cuts.

Proposition 15.3.3. *For any rational q, we have $\{r \in \mathbb{Q} \mid r < q\} \in \mathbb{R}$.*

PROOF:
$\langle 1 \rangle 1$. LET: $q \in \mathbb{Q}$
$\langle 1 \rangle 2$. LET: $q \downarrow = \{r \mid r < q\}$
$\langle 1 \rangle 3$. $q \downarrow \neq \emptyset$
    PROOF: We have $q - 1 \in q \downarrow$.
$\langle 1 \rangle 4$. $q \downarrow \neq \mathbb{Q}$
    PROOF: Since $q \notin q \downarrow$.
$\langle 1 \rangle 5$. $q \downarrow$ is closed downwards.
    PROOF: Trivial.
$\langle 1 \rangle 6$. $q \downarrow$ has no greatest element.
    PROOF: For all $r \in q \downarrow$ we have $r < (q + r)/2 \in q \downarrow$.
▯

Proposition 15.3.4. *For rationals q and r, we have $q = r$ if and only if $\{s \in \mathbb{Q} \mid s < q\} = \{s \in \mathbb{Q} \mid s < r\}$.*

PROOF:
$\langle 1 \rangle 1$. LET: $q \downarrow = \{s \in \mathbb{Q} \mid s < q\}$
$\langle 1 \rangle 2$. LET: $r \downarrow = \{s \in \mathbb{Q} \mid s < r\}$
$\langle 1 \rangle 3$. If $q = r$ then $q \downarrow = r \downarrow$
    PROOF: Trivial.

⟨1⟩4. If $q < r$ then $q \downarrow \neq r \downarrow$
    PROOF: We have $q \in r \downarrow$ and $q \notin q \downarrow$.
⟨1⟩5. If $r < q$ then $q \downarrow \neq r \downarrow$
    PROOF: We have $r \in q \downarrow$ and $q \notin q \downarrow$.
☐

    Henceforth we identify a rational q with the real number $\{r \in \mathbb{Q} \mid r < q\}$.

Definition 15.3.5. Define the ordering $<$ on $\mathbb{R}$ by: $x < y$ iff $x \subsetneq y$.

Proposition 15.3.6. *The ordering on the reals agrees with the ordering on the rationals.*

PROOF:
⟨1⟩1. LET: $q, r \in \mathbb{Q}$
⟨1⟩2. LET: $q \downarrow = \{s \in \mathbb{Q} \mid s < q\}$.
⟨1⟩3. LET: $r \downarrow = \{s \in \mathbb{Q} \mid s < r\}$.
      PROVE: $q < r$ iff $q \downarrow \subsetneq r \downarrow$
⟨1⟩4. If $q < r$ then $q \downarrow \subsetneq r \downarrow$
  ⟨2⟩1. ASSUME: $q < r$
  ⟨2⟩2. $q \downarrow \subseteq r \downarrow$
    PROOF: If $s < q$ then $s < r$.
  ⟨2⟩3. $q \downarrow \neq r \downarrow$
    PROOF: Proposition 15.3.4.
⟨1⟩5. If $q \downarrow \subsetneq r \downarrow$ then $q < r$
  ⟨2⟩1. ASSUME: $q \downarrow \subsetneq r \downarrow$
  ⟨2⟩2. PICK $s \in r \downarrow$ such that $s \notin q \downarrow$
  ⟨2⟩3. $q \leq s < r$
☐

Proposition 15.3.7. *The ordering $<$ is a linear ordering on $\mathbb{R}$.*

PROOF:
⟨1⟩1. $<$ is irreflexive.
    PROOF: No set is a proper subset of itself.
⟨1⟩2. $<$ is transitive.
    PROOF: Since the relationship $\subsetneq$ is transitive on the class of all sets.
⟨1⟩3. $<$ is total.
  ⟨2⟩1. LET: x, y be Dedekind cuts.
  ⟨2⟩2. ASSUME: $x \nsubseteq y$
      PROVE: $y \subsetneq x$
  ⟨2⟩3. PICK $q \in x$ such that $q \notin y$
  ⟨2⟩4. LET: $r \in y$
      PROVE: $r \in x$
  ⟨2⟩5. $q \nleq r$
    PROOF: Since y is closed downwards.
  ⟨2⟩6. $r < q$
  ⟨2⟩7. $r \in x$

PROOF: Since x is closed downwards.

□

Proposition 15.3.8. *Any bounded nonempty subset of* $\mathbb{R}$ *has a least upper bound.*

PROOF:
⟨1⟩1. LET: A be a bounded nonempty subset of $\mathbb{R}$.
⟨1⟩2. $\bigcup A$ is a Dedekind cut.
  ⟨2⟩1. $\bigcup A \neq \emptyset$
    ⟨3⟩1. PICK $x \in A$
    ⟨3⟩2. PICK $q \in x$
    ⟨3⟩3. $q \in \bigcup A$
  ⟨2⟩2. $\bigcup A \neq \mathbb{Q}$
    ⟨3⟩1. PICK an upper bound u for A
    ⟨3⟩2. PICK $q \notin u$
      PROVE: $q \notin \bigcup A$
    ⟨3⟩3. ASSUME: for a contradiction $q \in \bigcup A$
    ⟨3⟩4. PICK $x \in A$ such that $q \in x$
    ⟨3⟩5. $x \leq u$
    ⟨3⟩6. $q \in u$
    ⟨3⟩7. Q.E.D.
      PROOF: This is a contradiction.
  ⟨2⟩3. $\bigcup A$ is closed downwards.
    ⟨3⟩1. LET: $q \in \bigcup A$ and $r < q$
    ⟨3⟩2. PICK $x \in A$ such that $q \in x$
    ⟨3⟩3. $r \in x$
    ⟨3⟩4. $r \in \bigcup A$
  ⟨2⟩4. $\bigcup A$ has no greatest element.
    ⟨3⟩1. LET: $q \in \bigcup A$
    ⟨3⟩2. PICK $x \in A$ such that $q \in x$
    ⟨3⟩3. PICK $r \in x$ such that $q < r$
    ⟨3⟩4. $r \in \bigcup A$
⟨1⟩3. $\bigcup A$ is an upper bound for A.
  PROOF: For all $x \in A$ we have $x \subseteq \bigcup A$.
⟨1⟩4. For any upper bound u for $\bigcup A$ we have $\bigcup A \leq u$.
  PROOF: If $\forall x \in A.x \subseteq u$ we have $\bigcup A \subseteq u$.
□

Definition 15.3.9 (Addition)**.** Define *addition* $+$ on the reals by

$$x + y := \{q + r \mid q \in x, r \in y\} \ .$$

We prove this is well-defined.

PROOF:
⟨1⟩1. LET: $x, y \in \mathbb{R}$
    PROVE: $X + y$ is a Dedekind cut.

$\langle 1 \rangle$2. $x + y \neq \emptyset$
   PROOF: Pick $q \in x$ and $r \in y$; then $q + r \in x + y$.
$\langle 1 \rangle$3. $x + y \neq \mathbb{Q}$
   $\langle 2 \rangle$1. PICK $q \notin x$ and $r \notin y$
      PROVE: $q + r \notin x + y$
   $\langle 2 \rangle$2. ASSUME: for a contradiction $q + r \in x + y$
   $\langle 2 \rangle$3. PICK $q' \in x$ and $r' \in y$ such that $q + r = q' + r'$
   $\langle 2 \rangle$4. $q' < q$ and $r' < r$
   $\langle 2 \rangle$5. $q' + r' < q + r$
   $\langle 2 \rangle$6. Q.E.D.
     PROOF: This is a contradiction.
$\langle 1 \rangle$4. $x + y$ is closed downwards.
   $\langle 2 \rangle$1. LET: $q \in x$ and $r \in y$
   $\langle 2 \rangle$2. LET: $s < q + r$
      PROVE: $s \in x + y$
   $\langle 2 \rangle$3. $s - r < q$
   $\langle 2 \rangle$4. $s - r \in x$
   $\langle 2 \rangle$5. $s = (s - r) + r \in x + y$
$\langle 1 \rangle$5. $x + y$ has no greatest element.
   $\langle 2 \rangle$1. LET: $q \in x$ and $r \in y$
      PROVE: There exists $s \in x + y$ such that $q + r < s$
   $\langle 2 \rangle$2. PICK $q' \in x$ and $r' \in y$ such that $q < q'$ and $r < r'$
   $\langle 2 \rangle$3. $q + r < q' + r' \in x + y$
☐

Proposition 15.3.10. *Addition on the reals agrees with addition on the rationals.*

PROOF:
$\langle 1 \rangle$1. LET: $q, r \in \mathbb{Q}$
$\langle 1 \rangle$2. $q \downarrow + r \downarrow \subseteq (q + r) \downarrow$
   PROOF: If $s_1 < q$ and $s_2 < r$ then $s_1 + s_2 < q + r$.
$\langle 1 \rangle$3. $(q + r) \downarrow \subseteq q \downarrow + r \downarrow$
   $\langle 2 \rangle$1. LET: $s < q + r$
   $\langle 2 \rangle$2. $s - r < q$
   $\langle 2 \rangle$3. PICK t such that $s - r < t < q$
   $\langle 2 \rangle$4. $s - t < r$
   $\langle 2 \rangle$5. $s = t + (s - t) \in q \downarrow + r \downarrow$
☐

Proposition 15.3.11. *Addition is associative.*

PROOF:
$$x + (y + z) = \{q + r \mid q \in x, r \in y + z\}$$
$$= \{q + s_1 + s_2 \mid q \in x, s_1 \in y, s_2 \in z\}$$
$$= \{r + s_2 \mid r \in x + y, s_2 \in z\}$$
$$= (x + y) + z \qquad\qquad ☐$$

Proposition 15.3.12. *Addition is commutative.*

PROOF:
$$x + y = \{q + r \mid q \in x, r \in y\}$$
$$= \{r + q \mid r \in y, q \in x\}$$
$$= y + x \qquad\qquad\qquad \square$$

Proposition 15.3.13. *For any $x \in \mathbb{R}$ we have $x + 0 = x$.*

PROOF:
$\langle 1 \rangle 1.$ $x + 0 \subseteq x$
   PROOF: If $q \in x$ and $r < 0$ then $q + r < q$ so $q + r \in x$.
$\langle 1 \rangle 2.$ $x \subseteq x + 0$
   $\langle 2 \rangle 1.$ LET: $q \in x$
   $\langle 2 \rangle 2.$ PICK $r \in x$ such that $q < r$.
     PROOF: x has no greatest element.
   $\langle 2 \rangle 3.$ $q - r < 0$
   $\langle 2 \rangle 4.$ $q = r + (q - r) \in x + 0$
$\square$

Definition 15.3.14. For $x \in \mathbb{R}$, define $-x := \{q \in \mathbb{Q} \mid \exists r > q. -r \notin x\}$.

Proposition 15.3.15. *For all $x \in \mathbb{R}$ we have $-x \in \mathbb{R}$.*

PROOF:
$\langle 1 \rangle 1.$ LET: $x \in \mathbb{R}$
$\langle 1 \rangle 2.$ $-x \neq \emptyset$
   $\langle 2 \rangle 1.$ PICK $s \notin x$
   $\langle 2 \rangle 2.$ $-s - 1 \in -x$
$\langle 1 \rangle 3.$ $-x \neq \mathbb{Q}$
   $\langle 2 \rangle 1.$ PICK $s \in x$
      PROVE: $-s \notin -x$
   $\langle 2 \rangle 2.$ ASSUME: for a contradiction $-s \in -x$
   $\langle 2 \rangle 3.$ PICK $r > -s$ such that $-r \notin x$
   $\langle 2 \rangle 4.$ $-r < s$
   $\langle 2 \rangle 5.$ Q.E.D.
     PROOF: This contradicts the fact that x is closed downwards.
$\langle 1 \rangle 4.$ $-x$ is closed downwards.
   PROOF: Immediate from definition.
$\langle 1 \rangle 5.$ $-x$ has no greatest element.
   $\langle 2 \rangle 1.$ LET: $q \in -x$
   $\langle 2 \rangle 2.$ PICK $r > q$ such that $-r \notin x$
   $\langle 2 \rangle 3.$ PICK s such that $q < s < r$
   $\langle 2 \rangle 4.$ $s \in -x$
$\square$

Lemma 15.3.16. *Let p be a positive rational number. For any real number x, there exists a rational $q \in x$ such that $p + q \notin x$.*

PROOF:

$\langle 1 \rangle$1. PICK $q_0 \in x$

$\langle 1 \rangle$2. There exists $k \in \mathbb{N}$ such that $q_0 + kp \notin x$

   $\langle 2 \rangle$1. PICK $q_1 \notin x$

   $\langle 2 \rangle$2. PICK $k \in \mathbb{N}$ such that $q_1 - q_0 < pk$

     PROOF: Proposition 15.2.11.

   $\langle 2 \rangle$3. $q_1 < q_0 + kp$

   $\langle 2 \rangle$4. $q_0 + kp \notin x$

$\langle 1 \rangle$3. LET: k be the least natural number such that $q_0 + kp \notin x$

$\langle 1 \rangle$4. $k \neq 0$

   PROOF: $\langle 1 \rangle$1

$\langle 1 \rangle$5. LET: $q = q_0 + (k-1)p$

$\langle 1 \rangle$6. $q \in x$ and $q + p \notin x$.

$\square$

Proposition 15.3.17. *For every real x we have $x + (-x) = 0$.*

PROOF:

$\langle 1 \rangle$1. LET: x be a real number.

$\langle 1 \rangle$2. $x + (-x) \subseteq 0$

   $\langle 2 \rangle$1. LET: $q_1 \in x$ and $q_2 \in -x$

   $\langle 2 \rangle$2. PICK $r > q_2$ such that $-r \notin x$

   $\langle 2 \rangle$3. $q_1 < -r$

   $\langle 2 \rangle$4. $r < -q_1$

   $\langle 2 \rangle$5. $q_2 < -q_1$

   $\langle 2 \rangle$6. $q_1 + q_2 < 0$

$\langle 1 \rangle$3. $0 \subseteq x + (-x)$

   $\langle 2 \rangle$1. LET: $p < 0$

   $\langle 2 \rangle$2. $0 < -p$

   $\langle 2 \rangle$3. PICK $q \in x$ such that $q - p/2 \notin x$

     PROOF: Lemma 15.3.16.

   $\langle 2 \rangle$4. LET: $s = p/2 - q$

   $\langle 2 \rangle$5. $-s \notin x$

   $\langle 2 \rangle$6. $p - q < s$

   $\langle 2 \rangle$7. $p - q \in -x$

   $\langle 2 \rangle$8. $p \in x + (-x)$

$\square$

Corollary 15.3.17.1. *The reals form an Abelian group under addition.*

Proposition 15.3.18. *For any reals x, y and z, we have $x < y$ if and only if $x + z < y + z$.*

PROOF:

$\langle 1 \rangle$1. $\forall x, y, z \in \mathbb{R}.x \leq y \Rightarrow x + z \leq y + z$

   $\langle 2 \rangle$1. LET: $x, y, z \in \mathbb{R}$

   $\langle 2 \rangle$2. ASSUME: $x \leq y$

   $\langle 2 \rangle$3. For all $q \in x$ and $r \in z$ we have $q + r \in y + z$

$\langle 1\rangle 2.$ $\forall x, y, z \in \mathbb{R}.x + z = y + z \Leftrightarrow x = y$
　　PROOF: Proposition 12.1.4.
$\langle 1\rangle 3.$ $\forall x, y, z \in \mathbb{R}.x < y \Rightarrow x + z < y + z$
$\langle 1\rangle 4.$ Q.E.D.
　　PROOF: Proposition 5.2.6.
▯

Definition 15.3.19 (Absolute Value)**.** The *absolute value* of a real number x is defined to be
$$|x| = \begin{cases} x & \text{if } 0 \le x \\ -x & \text{if } x < 0 \end{cases}$$

Definition 15.3.20 (Multiplication)**.** Define *multiplication* $\cdot$ on $\mathbb{R}$ as follows:

- If x and y are non-negative then

$$xy = 0 \cup \{rs \mid 0 \le r \in x \wedge 0 \le s \in y\} \ .$$

- If x and y are both negative then

$$xy = (-x)(-y) \ .$$

- If one of x and y is negative and one is non-negative then

$$xy = -(|x||y|) \ .$$

We prove this is well-defined.

PROOF:
$\langle 1\rangle 1.$ LET: x and y be non-negative reals.
　　　PROVE:　xy is real.
$\langle 1\rangle 2.$ $xy \ne \emptyset$
　　PROOF: Since $-1 \in xy$.
$\langle 1\rangle 3.$ $xy \ne \mathbb{Q}$
　　$\langle 2\rangle 1.$ PICK $r \notin x$ and $s \notin y$
　　　　PROVE:　$rs \notin xy$
　　$\langle 2\rangle 2.$ $0 \le r$ and $0 \le s$
　　　PROOF: Since $0 \subseteq x$ and $0 \subseteq y$.
　　$\langle 2\rangle 3.$ ASSUME: for a contradiction $rs \in xy$
　　$\langle 2\rangle 4.$ PICK r' and s' such that $0 \le r' \in x$, $0 \le s' \in y$ and $rs = r's'$
　　$\langle 2\rangle 5.$ $r' < r$
　　$\langle 2\rangle 6.$ $s' < s$
　　$\langle 2\rangle 7.$ $r's' < rs$
　　$\langle 2\rangle 8.$ Q.E.D.
　　　PROOF: This is a contradiction.
$\langle 1\rangle 4.$ xy is closed downwards.
　　$\langle 2\rangle 1.$ LET: $q \in xy$ and $r < q$

$\langle 2 \rangle 2$. CASE: $q \in 0$
  PROOF: Then $r < q < 0$ so $r \in xy$
$\langle 2 \rangle 3$. CASE: $q = s_1 s_2$ where $0 \le s_1 \in x$ and $0 \le s_2 \in y$
  $\langle 3 \rangle 1$. ASSUME: w.l.o.g. $0 \le r$
  $\langle 3 \rangle 2$. $0 < s_1$ and $0 < s_2$
  $\langle 3 \rangle 3$. $r/s_2 < s_1$
  $\langle 3 \rangle 4$. $r/s_2 \in x$
  $\langle 3 \rangle 5$. $r = (r/s_2)s_2 \in xy$
$\langle 1 \rangle 5$. xy has no greatest element.
  $\langle 2 \rangle 1$. LET: $q \in xy$
  $\langle 2 \rangle 2$. CASE: $q \in 0$
    PROOF: $q < q/2 \in 0$
  $\langle 2 \rangle 3$. CASE: $q = rs$ where $0 \le r \in x$ and $0 \le s \in y$
    $\langle 3 \rangle 1$. PICK r' and s' with $r < r' \in x$ and $s < s' \in y$
    $\langle 3 \rangle 2$. $q < r's' \in xy$
$\square$

Proposition 15.3.21. *Multiplication is commutative.*

PROOF: Immediate from definition. $\square$

Proposition 15.3.22. *Multiplication is associative.*

PROOF:
$\langle 1 \rangle 1$. For non-negative reals x, y and z, we have $x(yz) = (xy)z$
  PROOF: It computes to $0 \cup \{qrs \mid 0 \le q \in x, 0 \le r \in y, 0 \le s \in z\}$.
$\langle 1 \rangle 2$. For all reals x, y and z, we have $x(yz) = (xy)z$
  PROOF: It is equal to $|x||y||z|$ if an even number of them are negative, and
  $-(|x||y||z|)$ otherwise.
$\square$

Proposition 15.3.23. *Multiplication distributes over addition.*

PROOF:
$\langle 1 \rangle 1$. For all non-negative reals x, y and z, we have $x(y + z) = xy + xz$
  $\langle 2 \rangle 1$. LET: x, y and z be non-negative reals.
  $\langle 2 \rangle 2$. $x(y + z) \subseteq xy + xz$
    $\langle 3 \rangle 1$. LET: $q \in x(y + z)$
    $\langle 3 \rangle 2$. CASE: $q < 0$
      PROOF: Then we have $q/2 \in xy$ and $q/2 \in xz$ so $q \in xy + xz$.
    $\langle 3 \rangle 3$. CASE: $q = rs$ where $0 \le r \in x$ and $0 \le s \in y + z$
      $\langle 4 \rangle 1$. PICK $s_1 \in y$ and $s_2 \in z$ such that $s = s_1 + s_2$
      $\langle 4 \rangle 2$. $rs_1 \in xy$
        PROOF: If $s_1 < 0$ then $rs_1 < 0$ so $rs_1 \in xy$. If $0 \le s_1$ then we also
        have $rs_1 \in xy$.
      $\langle 4 \rangle 3$. $rs_2 \in xz$
        PROOF: Similar.
      $\langle 4 \rangle 4$. $q \in xy + xz$

PROOF: Since $q = rs_1 + rs_2$.

$\langle 2 \rangle 3.$ $xy + xz \subseteq x(y + z)$

    $\langle 3 \rangle 1.$ LET: $q \in xy$ and $r \in xz$.

        PROVE: $q + r \in x(y + z)$

    $\langle 3 \rangle 2.$ CASE: $q < 0$ and $r < 0$

      PROOF: Then $q + r < 0$ so $q + r \in x(y + z)$.

    $\langle 3 \rangle 3.$ CASE: $q < 0$ and $r = r_1 r_2$ where $0 \leq r_1 \in x$ and $0 \leq r_2 \in z$

        $\langle 4 \rangle 1.$ $q + r < r$

        $\langle 4 \rangle 2.$ $q + r \in xz$

        $\langle 4 \rangle 3.$ ASSUME: w.l.o.g. $0 \leq q + r$

          PROOF: Otherwise $q + r \in x(y + z)$ immediately.

        $\langle 4 \rangle 4.$ PICK s_1, s_2 with $0 \leq s_1 \in x$, $0 \leq s_2 \in y$ and $q + r = s_1 s_2$

        $\langle 4 \rangle 5.$ $s_2 \in y + z$

          PROOF: Since $0 \in z$ so $s_2 = s_2 + 0 \in y + z$.

        $\langle 4 \rangle 6.$ $q + r \in x(y + z)$

    $\langle 3 \rangle 4.$ CASE: $q = q_1 q_2$ where $0 \leq q_1 \in x$ and $0 \leq q_2 \in y$ and $r < 0$

      PROOF: Similar.

    $\langle 3 \rangle 5.$ CASE: $q = q_1 q_2$ where $0 \leq q_1 \in x$ and $0 \leq q_2 \in y$ and $r = r_1 r_2$ where
              $0 \leq r_1 \in x$ and $0 \leq r_2 \in z$

        $\langle 4 \rangle 1.$ ASSUME: w.l.o.g. $q_1 \leq r_1$

        $\langle 4 \rangle 2.$ $q + r \leq r_1(q_2 + r_2) \in x(y + z)$

$\langle 1 \rangle 2.$ For any negative real x and non-negative reals y and z, we have $x(y+z) = xy + xz$

  PROOF:

$$x(y + z) = -(-x)(y + z) = -((-x)y + (-x)z) \qquad (\langle 1 \rangle 1)$$
$$= -((-x)y) - ((-x)z)$$
$$= xy + xz$$

$\langle 1 \rangle 3.$ For any non-negative real x and reals y and z with one negative and one
      non-negative, we have $x(y + z) = xy + xz$

  $\langle 2 \rangle 1.$ ASSUME: w.l.o.g. y is negative and z is non-negative.

  $\langle 2 \rangle 2.$ CASE: $0 \leq y + z$

    PROOF:

$$xy + xz = xy + x(-y + y + z)$$
$$= -(x(-y)) + x(-y + y + z)$$
$$= -(x(-y)) + x(-y) + x(y + z) \qquad (\langle 1 \rangle 1)$$
$$= x(y + z)$$

  $\langle 2 \rangle 3.$ CASE: $y + z < 0$

    $\langle 3 \rangle 1.$ $-y - z > 0$

    $\langle 3 \rangle 2.$ $-y = z - y - z$

    $\langle 3 \rangle 3.$ $xy + xz = x(y + z)$

PROOF:

$$\begin{aligned}
xy + xz &= -(x(-y)) + xz \\
&= -(x(z - y - z)) + xz \\
&= -(xz + x(-y - z)) + xz \qquad (\langle 1 \rangle 1) \\
&= -xy - x(-y - z) + xz \\
&= -x(-y - z) \\
&= x(y + z)
\end{aligned}$$

$\langle 1 \rangle 4$. For any non-negative real x and negative reals y and z, we have $x(y+z) = xy + xz$

PROOF:

$$\begin{aligned}
x(y + z) &= -x(-y - z) \\
&= -(x(-y) + x(-z)) \qquad (\langle 1 \rangle 1) \\
&= -x(-y) - x(-z) \\
&= xy + xz
\end{aligned}$$

$\langle 1 \rangle 5$. For any negative real x and reals y and z with one negative and one non-negative, we have $x(y + z) = xy + xz$

$\langle 2 \rangle 1$. ASSUME: w.l.o.g. y is negative and z is non-negative.

$\langle 2 \rangle 2$. CASE: $0 \le y + z$

PROOF:

$$\begin{aligned}
x(y + z) &= -((-x)(y + z)) \\
&= -((-x)y + (-x)z) \qquad (\langle 1 \rangle 3) \\
&= -((-x)y) - ((-x)z) \\
&= (-x)(-y) - ((-x)z) \\
&= xy + xz
\end{aligned}$$

$\langle 2 \rangle 3$. CASE: $y + z < 0$

PROOF:

$$\begin{aligned}
x(y + z) &= (-x)(-y - z) \\
&= (-x)(-y) + (-x)(-z) \qquad (\langle 1 \rangle 3) \\
&= xy + xz
\end{aligned}$$

$\langle 1 \rangle 6$. For any negative reals x, y and z, we have $x(y + z) = xy + xz$

PROOF:

$$\begin{aligned}
x(y + z) &= (-x)(-y - z) \\
&= (-x)(-y) + (-x)(-z) \qquad (\langle 1 \rangle 1) \\
&= xy + xz
\end{aligned}$$

▯

Proposition 15.3.24. *For any real x we have $x1 = x$.*

PROOF:

$\langle 1 \rangle 1$. CASE: $0 \le x$

$\quad \langle 2 \rangle 1$. $x1 \subseteq x$

$\quad \quad \langle 3 \rangle 1$. LET: $q \in x1$

$\langle 3 \rangle 2$. CASE: $q < 0$
    PROOF: Then $q \in x$ because $0 \leq x$.
  $\langle 3 \rangle 3$. $q = rs$ where $0 \leq r \in x$ and $0 \leq s < 1$
    PROOF: Then $q < r$ so $q \in x$.
$\langle 2 \rangle 2$. $x \subseteq x1$
  $\langle 3 \rangle 1$. LET: $q \in x$
  $\langle 3 \rangle 2$. ASSUME: w.l.o.g. $0 \leq q$
  $\langle 3 \rangle 3$. PICK r such that $q < r \in x$
  $\langle 3 \rangle 4$. $0 \leq q/r < 1$
  $\langle 3 \rangle 5$. $q = r(q/r) \in x1$
$\langle 1 \rangle 2$. CASE: $x < 0$
  PROOF:

$$x1 = -((-x)1)$$
$$= -(-x) \qquad\qquad (\langle 1 \rangle 1)$$
$$= x$$

□

Lemma 15.3.25. *Let $x \in \mathbb{R}$ and c be a positive rational. Then there exists $a \in x$ and a non-least rational upper bound b for x such that $b - a = c$.*

PROOF:
$\langle 1 \rangle 1$. PICK $a_1 \in x$ such that if x has a rational supremum s then $a_1 > s - c$
$\langle 1 \rangle 2$. There exists a natural number n such that $a_1 + nc$ is an upper bound for
    x.
  $\langle 2 \rangle 1$. PICK a non-least upper bound b_1 for x.
  $\langle 2 \rangle 2$. PICK a natural number n such that $nc > b_1 - a_1$
    PROOF: Proposition 15.2.11.
  $\langle 2 \rangle 3$. $a_1 + nc > b_1$
  $\langle 2 \rangle 4$. $a_1 + nc$ is an upper bound for x.
$\langle 1 \rangle 3$. LET: k be the least natural number such that $a_1 + kc$ is an upper bound
      for x.
$\langle 1 \rangle 4$. $a_1 + (k - 1)c \in x$
$\langle 1 \rangle 5$. $a_1 + kc$ is not the supremum of x.
  $\langle 2 \rangle 1$. ASSUME: for a contradiction $a_1 + kc$ is the supremum of x.
  $\langle 2 \rangle 2$. $a_1 > a_1 + (k - 1)c$
    PROOF: $\langle 1 \rangle 1$
  $\langle 2 \rangle 3$. Q.E.D.
    PROOF: This is a contradiction.
$\langle 1 \rangle 6$. LET: $a = a_1 + (k - 1)c$
$\langle 1 \rangle 7$. LET: $b = a_1 + kc$
$\langle 1 \rangle 8$. $b - a = c$
□

Proposition 15.3.26. *For any non-zero real x, there exists a real y such that $xy = 1$.*

PROOF:

⟨1⟩1. CASE: $0 < x$

  ⟨2⟩1. LET: $y = \{q \in \mathbb{Q} \mid q \le 0\} \cup \{u^{-1} \mid u$ is an upper bound for x but not the supremum of $x\}$

  ⟨2⟩2. y is a real number.

    ⟨3⟩1. $y \ne \emptyset$

      PROOF: Since $0 \in y$.

    ⟨3⟩2. $y \ne \mathbb{Q}$

      ⟨4⟩1. PICK $q \in x$ such that $0 < q$

      ⟨4⟩2. $q^{-1} \notin y$

    ⟨3⟩3. y is closed downwards.

      ⟨4⟩1. LET: $q \in y$ and $r < q$

        PROVE: $r \in y$

      ⟨4⟩2. ASSUME: w.l.o.g. $0 < r$

      ⟨4⟩3. q^{-1} is a non-least upper bound for x.

      ⟨4⟩4. $q^{-1} < r^{-1}$

      ⟨4⟩5. r^{-1} is a non-least upper bound for x.

      ⟨4⟩6. $r \in y$

    ⟨3⟩4. y has no greatest element.

      ⟨4⟩1. LET: $q \in y$

        PROVE: There exists $r \in y$ such that $q < r$

      ⟨4⟩2. CASE: $q \le 0$

        ⟨5⟩1. PICK a non-least upper bound u for x.

        ⟨5⟩2. $q < u^{-1} \in x$

      ⟨4⟩3. CASE: $q = u^{-1}$ where u is a non-least upper bound for x.

        ⟨5⟩1. PICK a non-least upper bound v with $v < u$

        ⟨5⟩2. $u^{-1} < v^{-1} \in y$

  ⟨2⟩3. $0 < y$

  ⟨2⟩4. $xy \subseteq 1$

    ⟨3⟩1. LET: $q \in xy$

    ⟨3⟩2. ASSUME: w.l.o.g. $0 < q$

    ⟨3⟩3. PICK $0 < r \in x$ and $0 < s \in y$ such that $q = rs$

    ⟨3⟩4. s^{-1} is a non-least upper bound for x

    ⟨3⟩5. $r < s^{-1}$

    ⟨3⟩6. $rs < 1$

  ⟨2⟩5. $1 \subseteq xy$

    ⟨3⟩1. LET: $q < 1$

      PROVE: $q \in xy$

    ⟨3⟩2. ASSUME: w.l.o.g. $0 < q$

    ⟨3⟩3. PICK a_1 with $0 < a_1 \in x$

    ⟨3⟩4. $(1-q)a_1 > 0$

    ⟨3⟩5. PICK $a \in x$ and a non-least upper bound w of x such that $w - a = (1-q)a_1$

      PROOF: Lemma 15.3.25.

    ⟨3⟩6. $w - a < (1-q)w$

    ⟨3⟩7. $qw < a$

    ⟨3⟩8. $w < a/q$

    ⟨3⟩9. a/q is a non-least upper bound for x

    $\langle 3 \rangle$10. $q/a \in y$
    $\langle 3 \rangle$11. $q \in xy$
$\langle 1 \rangle$2. CASE: $x < 0$
  $\langle 2 \rangle$1. PICK y such that $(-x)y = 1$
    PROOF: $\langle 1 \rangle$1
  $\langle 2 \rangle$2. $x(-y) = 1$
☐

Proposition 15.3.27. *For real numbers x, y and z, if $0 < z$ then $x < y$ if and only if $xz < yz$.*

PROOF:
$\langle 1 \rangle$1. For any real numbers x, y and z, if $0 < z$ and $x < y$ then $xz < yz$
  $\langle 2 \rangle$1. LET: x, y and z be real numbers.
  $\langle 2 \rangle$2. ASSUME: $0 < z$ and $x < y$.
  $\langle 2 \rangle$3. $y = x + (y - x)$
  $\langle 2 \rangle$4. $y - x > 0$
  $\langle 2 \rangle$5. $(y - x)z > 0$
  $\langle 2 \rangle$6. $yz > xz$
    PROOF:

$$\begin{aligned} yz &= (x + (y - x))z \\ &= xz + (y - x)z \\ &> xz \end{aligned}$$

$\langle 1 \rangle$2. For any real numbers x, y and z, if $0 < z$ and $xz < yz$ then $x < y$
  PROOF: Proposition 5.2.6.
☐

Corollary 15.3.27.1. *The real numbers form a complete ordered field.*

Proposition 15.3.28.
$$(0, 1) \approx \mathbb{R}$$

PROOF: The function $f(x) = (2x - 1)/(x - x^2)$ is a bijection between $(0, 1)$ and $\mathbb{R}$. ☐

Proposition 15.3.29.
$$|\mathbb{R}| = 2^{\aleph_0}$$

PROOF:
$\langle 1 \rangle$1. $(0, 1) \preccurlyeq 2^{\mathbb{N}}$
  PROOF: The function H where $H(x)(n)$ is the nth binary digit of the binary expansion of x is an injection.
$\langle 1 \rangle$2. $2^{\mathbb{N}} \preccurlyeq \mathbb{R}$
  PROOF: Map f to the real number in $[0, 1/9]$ whose $n + 1$st decimal digit is $f(n)$.
☐

Proposition 15.3.30. *The set of algebraic numbers is countable.*

PROOF: There are countably many integer polynomials, each with finitely many roots. ☐

Corollary 15.3.30.1. *There are uncountably many transcendental numbers.*

Proposition 15.3.31. *Let A be a set of disks in the plane, no two of which intersect. Then A is countable.*

PROOF: Every circle includes a point with rational coordinates. Define $f : \{q \in \mathbb{Q}^2 \mid \exists C \in A.q \in C\} \to A$ by $f(q) = C$ iff $q \in C$. Then f is surjective. ☐

Proposition 15.3.32. *There exists an uncountable set of circles in the plane that do not intersect.*

PROOF: The set of all circles with origin O is uncountable. ☐

Chapter 16

Real Analysis

Theorem 16.0.1 (Weierstrass)**.** *Let $a, b \in \mathbb{R}$ with $a < b$. Let $f : [a, b] \to \mathbb{R}$ be continuous. For every $\epsilon > 0$, there exists a polynomial p such that $\forall x \in [a, b].|f(x) - p(x)| < \epsilon$.*

Theorem 16.0.2 (Bolzano-Weierstrass)**.** *Every bounded sequence in $\mathbb{R}^n$ has a convergent subsequence.*

Chapter 17

Complex Analysis

Theorem 17.0.1 (Hölder's Inequality). *Let p and q be real numbers with $p > 1$,*
$q > 1$ and $1/p + 1/q = 1$. If $(x_n) \in l^p$ and $(y_n) \in l^q$ then

$$\sum_{n=0}^{\infty} |x_n y_n| \leq \left(\sum_{n=0}^{\infty} |x_n|^p \right)^{1/p} \left(\sum_{n=0}^{\infty} |y_n|^q \right)^{1/q}$$

PROOF:

$\langle 1 \rangle 1$. LET: p and q be real numbers with $p > 1$ and $q > 1$

$\langle 1 \rangle 2$. ASSUME: $1/p + 1/q = 1$

$\langle 1 \rangle 3$. LET: $(x_n) \in l^p$

$\langle 1 \rangle 4$. LET: $(y_n) \in l^q$

$\langle 1 \rangle 5$. ASSUME: w.l.o.g. $x_0 \neq 0$ and $y_0 \neq 0$

$\langle 1 \rangle 6$. For all $x \in [0, 1]$, we have

$$x^{1/p} \leq \frac{1}{p} x + \frac{1}{q} \quad .$$

   PROOF:

   $\langle 2 \rangle 1$. LET: $f : [0, 1] \to \mathbb{R}$ be the function

$$f(x) = \frac{1}{p} x + \frac{1}{q} - x^{1/p} \quad .$$

   $\langle 2 \rangle 2$. $f'(x) = \frac{1}{p} - \frac{1}{p} x^{-1/q}$ for $x \in (0, 1]$

   $\langle 2 \rangle 3$. $f'(x) < 0$ for $x \in (0, 1]$

   $\langle 2 \rangle 4$. $f(1) = 1/p + 1/q - 1 = 0$

   $\langle 2 \rangle 5$. $f(x) \geq 0$ for all $x \in [0, 1]$

$\langle 1 \rangle 7$. For all non-negative reals a and b, we have

$$ab \leq \frac{a^p}{p} + \frac{b^q}{q} \quad .$$

   $\langle 2 \rangle 1$. LET: a and b be non-negative reals.

   $\langle 2 \rangle 2$. CASE: $a^p \leq b^q$

      $\langle 3 \rangle 1$. $0 \leq a^p / b^q \leq 1$

      $\langle 3 \rangle 2$.

$$ab^{-q/p} \leq \frac{1}{p} \frac{a^p}{b^q} + \frac{1}{q}$$

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PROOF: Taking $x = a^p/b^q$ in $\langle 1 \rangle 6$.

$\langle 3 \rangle 3$.
$$ab^{1-q} \leq \frac{1}{p}\frac{a^p}{b^q} + \frac{1}{q}$$

PROOF: $-q/p = 1 - q$ from $\langle 1 \rangle 2$.

$\langle 3 \rangle 4$.
$$ab \leq \frac{a^p}{p} + \frac{b^q}{q}$$

$\langle 2 \rangle 3$. CASE: $b^q \leq a^p$

PROOF: Similar.

$\langle 1 \rangle 8$. For $j = 1, \ldots, n$, we have
$$\frac{|x_j|}{\left(\sum_{k=0}^{n}|x_k|^p\right)^{1/p}}\frac{|y_j|}{\left(\sum_{k=0}^{n}|y_k|^q\right)^{1/q}} \leq \frac{1}{p}\frac{|x_j|^p}{\sum_{k=0}^{n}|x_k|^p} + \frac{1}{q}\frac{|y_j|^q}{\sum_{k=0}^{n}|y_k|^q}$$

PROOF: From $\langle 1 \rangle 7$ with
$$a = \frac{|x_j|}{\left(\sum_{k=0}^{n}|x_k|^p\right)^{1/p}} \text{ and } b = \frac{|y_j|}{\left(\sum_{k=0}^{n}|y_k|^q\right)^{1/q}} \ .$$

$\langle 1 \rangle 9$.
$$\frac{\sum_{j=0}^{n}|x_j||y_j|}{\left(\sum_{k=0}^{n}|x_k|^p\right)^{1/p}\left(\sum_{k=0}^{n}|y_k|^q\right)^{1/q}} \leq 1$$

PROOF:
$$\frac{\sum_{j=0}^{n}|x_j||y_j|}{\left(\sum_{k=0}^{n}|x_k|^p\right)^{1/p}\left(\sum_{k=0}^{n}|y_k|^q\right)^{1/q}} \leq \frac{1}{p} + \frac{1}{q} \quad \text{(Taking the sum } j = 0 \text{ to } n \text{ in } \langle 1 \rangle 8\text{)}$$
$$= 1 \qquad\qquad (\langle 1 \rangle 2)$$

$\langle 1 \rangle 10$. Q.E.D.

PROOF: Taking the limit $n \to \infty$ in $\langle 1 \rangle 9$.

▯

Theorem 17.0.2 (Minkowski's Inequality). *Let p be a real number, $p \geq 1$. Let $(x_n), (y_n) \in l^p$. Then*

$$\left(\sum_{n=1}^{\infty}|x_n + y_n|^p\right)^{1/p} \leq \left(\sum_{n=1}^{\infty}|x_n|^p\right)^{1/p} + \left(\sum_{n=1}^{\infty}|y_n|^p\right)^{1/p}$$

PROOF:

$\langle 1 \rangle 1$. LET: p be a real number with $p \geq 1$

$\langle 1 \rangle 2$. ASSUME: w.l.o.g. $p > 1$

PROOF: The case $p = 1$ is just the Triangle Inequality.

$\langle 1 \rangle 3$. LET: q be the real such that $1/p + 1/q = 1$

$\langle 1 \rangle 4$.
$$\sum_{n=0}^{\infty}|x_n + y_n|^p \leq \left(\sum_{n=0}^{\infty}|x_n|^p\right)^{1/p}\left(\sum_{n=0}^{\infty}|x_n + y_n|^{q(p-1)}\right)^{1/q}$$
$$+ \left(\sum_{n=0}^{\infty}|y_n|^p\right)^{1/p}\left(\sum_{n=0}^{\infty}|x_n + y_n|^{q(p-1)}\right)^{1/q}$$

PROOF:

$$\sum_{n=0}^{\infty} |x_n + y_n|^p = \sum_{n=0}^{\infty} |x_n + y_n||x_n + y_n|^{p-1}$$

$$\leq \sum_{n=0}^{\infty} |x_n||x_n + y_n|^{p-1} + \sum_{n=0}^{\infty} |y_n||x_n + y_n|^{p-1} \quad \text{(Triangle Inequality)}$$

$$\leq \left(\sum_{n=0}^{\infty} |x_n|^p\right)^{1/p} \left(\sum_{n=0}^{\infty} |x_n + y_n|^{q(p-1)}\right)^{1/q}$$

$$+ \left(\sum_{n=0}^{\infty} |y_n|^p\right)^{1/p} \left(\sum_{n=0}^{\infty} |x_n + y_n|^{q(p-1)}\right)^{1/q} \quad \text{(Hölder's Inequality)}$$

$\langle 1 \rangle 5.$

$$\sum_{n=0}^{\infty} |x_n + y_n|^p \leq \left\{ \left(\sum_{n=0}^{\infty} |x_n|^p\right)^{1/p} + \left(\sum_{n=0}^{\infty} |y_n|^p\right)^{1/p} \right\} \left(\sum_{n=0}^{\infty} |x_n + y_n|^p\right)^{1/q}$$

$\langle 1 \rangle 6.$ Q.E.D.

$\Box$

Chapter 18

Topology

18.1 Topological Spaces

Definition 18.1.1 (Topology)**.** Let X be a set. A *topology* on X is a set $\mathcal{T} \subseteq \mathcal{P}X$, whose elements are called *open sets*, such that:

- $X \in \mathcal{T}$

- $\forall \mathcal{U} \subseteq \mathcal{T}. \bigcup \mathcal{U} \in \mathcal{T}$

- $\forall U, V \in \mathcal{T}. U \cap V \in \mathcal{T}$

A *topological space* is a pair $(X, \mathcal{T})$ such that X is a set and $\mathcal{T}$ is a topology on X. We refer to the elements of X as *points*.

An *open neighbourhood* of a point x is an open set U such that $x \in U$. We write $\mathcal{T}_x$ for the set of all open neighbourhoods of x.

Definition 18.1.2 (Closed Set)**.** In a topological space X, a set C is *closed* iff $X - C$ is open.

Definition 18.1.3 (Discrete Topology)**.** The *discrete topology* on a set X is $\mathcal{P}X$.

Definition 18.1.4 (Indiscrete Topology)**.** The *indiscrete topology* or *trivial topology* on a set X is $\{\emptyset, X\}$.

Definition 18.1.5 (Finer, Coarser)**.** Let $\mathcal{T}$ and $\mathcal{T}'$ be topologies on the same set X. Then $\mathcal{T}$ is *finer*, *larger* or *stronger* than $\mathcal{T}'$, and $\mathcal{T}'$ is *coarser*, *smaller* or *weaker* than $\mathcal{T}$, iff $\mathcal{T}' \subseteq \mathcal{T}$.

Definition 18.1.6 (Basis)**.** Let X be a set. A *basis* for a topology on X is a set $\mathcal{B} \subseteq \mathcal{P}X$, whose elements we call *basic open neighbourhoods*, such that:

- $\bigcup \mathcal{B} = X$

- $\forall A, B \in \mathcal{B}. \forall x \in A \cap B. \exists C \in \mathcal{B}. x \in C \subseteq A \cap B$.

The topology *generated* by $\mathcal{B}$ is the coarsest topology that includes $\mathcal{B}$.

Proposition 18.1.7. *The topology generated by $\mathcal{B}$ is $\{U \in \mathcal{P}X \mid \forall x \in U.\exists B \in \mathcal{B}.x \in B \subseteq U\}$*

18.2   Continuous Functions

Definition 18.2.1 (Continuous)**.** Let X and Y be topological spaces. Let $f : X \to Y$. Then f is *continuous* iff, for every open set U in Y, the set $f^{-1}(U)$ is open in X.

Proposition 18.2.2. *For any topological space X, the identity function $\mathrm{id}_X : X \to X$ is continuous.*

Proposition 18.2.3. *If $f : X \to Y$ and $g : Y \to Z$ are continuous then $g \circ f : X \to Z$ is continuous.*

Proposition 18.2.4. *Let X and Y be topological spaces and $f : X \to Y$. Then f is continuous if and only if, for every closed set C in Y, we have $f^{-1}(C)$ is closed in X.*

Definition 18.2.5 (Continuous at a Point)**.** Let X and Y be topological spaces. Let $f : X \to Y$. Let $x \in X$. Then f is *continuous at x* iff, for every open neighbourhood V of $f(x)$, we have $f^{-1}(V)$ is open.

18.3   Convergence

Definition 18.3.1 (Convergence)**.** Let X be a topological space. Let (x_n) be a sequence in X and $l \in X$. Then (x_n) *converges* to l, $x_n \to l$ as $n \to \infty$, if and only if, for every open neighbourhood U of l, there exists N such that $\forall n \geq N.x_n \in U$.

18.4   Metric Spaces

Definition 18.4.1 (Metric)**.** Let X be a set. A *metric* on a set X is a function $d : X^2 \to \mathbb{R}$ such that:

- $\forall x, y \in X.d(x, y) \geq 0$

- $\forall x, y \in X.d(x, y) = d(y, x)$

- *Triangle Inequality* $\forall x, y, z \in X.d(x, y) + d(y, z) \geq d(x, z)$

- $\forall x, y \in X.d(x, y) = 0$ iff $x = y$.

A *metric space* is a pair (X, d) such that d is a metric on X.

Definition 18.4.2 (Open Ball)**.** In a metric space X, let $c \in X$ and $r > 0$. The *open ball* with *centre* c and *radius* r is

$$B(c, r) := \{x \in X \mid d(x, c) < r\} \ .$$

Proposition 18.4.3. *In a metric space, the set of open balls forms a basis for a topology.*

Definition 18.4.4 (Metric Topology)**.** Given a metric space X, the *metric topology* on X is the topology generated by the basis of open balls.

A topological space $(X, \mathcal{T})$ is *metrizable* iff there exists a metric d on X such that $\mathcal{T}$ is the metric topology induced by d.

We identify a metric space with this topological space.

Proposition 18.4.5. *If d is a metric on X and $Y \subseteq X$ then $d \upharpoonright Y^2$ is a metric on Y.*

We write just Y for the metric space $(Y, d \upharpoonright Y^2)$.

Proposition 18.4.6. *Let X and Y be metric spaces. Let $f : X \to Y$. Then f is continuous if and only if, for every sequence (x_n) in X and $l \in X$, if $x_n \to l$ as $n \to \infty$ then $f(x_n) \to f(l)$ as $n \to \infty$.*

PROOF:
$\langle 1 \rangle 1$. If f is continuous then, for every sequence (x_n) in X and $l \in X$, if $x_n \to l$ as $n \to \infty$, then $f(x_n) \to f(l)$ as $n \to \infty$.
  $\langle 2 \rangle 1$. ASSUME: f is continuous.
  $\langle 2 \rangle 2$. LET: (x_n) be a sequence in X.
  $\langle 2 \rangle 3$. LET: $l \in X$
  $\langle 2 \rangle 4$. ASSUME: $x_n \to l$ as $n \to \infty$.
  $\langle 2 \rangle 5$. LET: V be an open neighbourhood of $f(l)$
  $\langle 2 \rangle 6$. $f^{-1}(V)$ is an open neighbourhood of l.
  $\langle 2 \rangle 7$. PICK N such that $\forall n \geq N.x_n \in f^{-1}(V)$
  $\langle 2 \rangle 8$. $\forall n \geq N.f(x_n) \in V$
$\langle 1 \rangle 2$. If, for every sequence (x_n) in X and $l \in X$, if $x_n \to l$ as $n \to \infty$, then $f(x_n) \to f(l)$ as $n \to \infty$, then f is continuous.
  $\langle 2 \rangle 1$. ASSUME: f is not continuous.
  $\langle 2 \rangle 2$. PICK an open set V in Y such that $f^{-1}(V)$ is not open in X
  $\langle 2 \rangle 3$. PICK $l \in f^{-1}(V)$ such that, for all $\epsilon > 0$, $B(l, \epsilon) \nsubseteq f^{-1}(V)$.
  $\langle 2 \rangle 4$. For $n \in \mathbb{N}$, PICK $x_n \in B(l, 1/(n+1))$ such that $x_n \notin f^{-1}(V)$.
  $\langle 2 \rangle 5$. $x_n \to l$ as $n \to \infty$.
  $\langle 2 \rangle 6$. $f(x_n) \nrightarrow f(l)$ as $n \to \infty$
□

Chapter 19

Linear Algebra

19.1 Vector Spaces

Definition 19.1.1 (Vector Space). Let K be either $\mathbb{R}$ or $\mathbb{C}$. A *vector space* over K consists of:

- a set V, whose elements are called *vectors*;

- an operation $+ : V^2 \to V$, *addition*;

- an operation $\cdot : K \times V \to V$, *scalar multiplication*

such that:

- V is an Abelian group under $+$

- $\forall \alpha, \beta \in K.\forall x \in V.\alpha(\beta x) = (\alpha\beta)x$

- $\forall \alpha, \beta \in K.\forall x \in V.(\alpha + \beta)x = \alpha x + \beta x$

- $\forall \alpha \in K.\forall x, y \in V.\alpha(x + y) = \alpha x + \alpha y$

- $\forall x \in V.1x = x$

We call the elements of K *scalars*. A *real vector space* is a vector space over $\mathbb{R}$, and a *complex vector space* is a vector space over $\mathbb{C}$.

Proposition 19.1.2. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a vector space over K. For any $\lambda \in K$ we have $\lambda 0 = 0$.*

PROOF:
$$\lambda 0 = \lambda(0 + 0)$$
$$= \lambda 0 + \lambda 0$$
$$\therefore 0 = \lambda 0 \qquad \qquad \square$$

Proposition 19.1.3. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a vector space over K. Let $\lambda \in K$ and $x \in V$. If $\lambda x = 0$ then either $\lambda = 0$ or $x = 0$.*

PROOF: If $\lambda \neq 0$ then $x = 1x = \lambda^{-1}\lambda x = \lambda^{-1}0 = 0$. $\square$

Proposition 19.1.4. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a vector space over K. For any $x \in V$ we have $0x = 0$.*

PROOF:
$$0x = (0 + 0)x$$
$$= 0x + 0x$$
$$\therefore 0 = 0x \qquad\qquad\qquad \square$$

Proposition 19.1.5. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a vector space over K. For any $x \in V$, we have $(-1)x = -x$.*

PROOF:
$$x + (-1)x = 1x + (-1)x$$
$$= (1 + (-1))x$$
$$= 0x$$
$$= 0$$
$$\therefore (-1)x = -x \qquad\qquad\qquad \square$$

Proposition 19.1.6. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Then K is a vector space over K under addition and multiplication in K.*

PROOF: Easy. $\square$

Proposition 19.1.7. *$\mathbb{C}$ is a vector space over $\mathbb{R}$.*

PROOF: Easy. $\square$

Proposition 19.1.8. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let $\{V_i\}_{i \in I}$ be a family of vector spaces over K. Then $\prod_{i \in I} V_i$ is a vector space under*

$$(f + g)(i) = f(i) + g(i) \qquad\qquad \left(f, g \in \prod_{i \in I} V_i, x \in X\right)$$

$$(\lambda f)(x) = \lambda f(x) \qquad\qquad \left(\lambda \in K, f \in \prod_{i \in I} V_i, x \in X\right)$$

PROOF: Easy. $\square$

19.2   Subspaces

Definition 19.2.1 (Vector Subspace)**.** Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a vector space over K. A *vector subspace* of V is a subset $U \subseteq V$ such that, for all $\alpha, \beta \in K$ and $x, y \in U$, we have $\alpha x + \beta y \in U$.

   It is a *proper* subspace iff $U \neq V$.

Proposition 19.2.2. *If U is a subspace of V then U is a vector space under the restrictions of $+$ and $\cdot$ to U.*

PROOF: Easy. $\square$

Proposition 19.2.3. *V is a subspace of V.*

PROOF: Easy. $\square$

Proposition 19.2.4. *If U is a subspace of V and V is a subspace of W then U is a subspace of W.*

PROOF: Easy. $\square$

Definition 19.2.5. Let Ω be a topological space. Then $\mathcal{C}(\Omega)$ is the complex vector space of all continuous functions from Ω to $\mathbb{C}$. This is a subspace of $\mathbb{C}^\Omega$.

Definition 19.2.6. Let $n, k \in \mathbb{N}$. Let Ω be an open subset of $\mathbb{R}^n$. Then $\mathcal{C}^k(\Omega)$ is the complex vector space of all functions $\Omega \to \mathbb{C}$ that have all continuous partial derivatives of order k. This is a subspace of $\mathcal{C}(\Omega)$. If $l > k$ then $\mathcal{C}^l(\Omega)$ is a subpase of $\mathcal{C}^k(\Omega)$.

Definition 19.2.7. Let $n \in \mathbb{N}$. Let Ω be an open subset of $\mathbb{R}^n$. Then $\mathcal{C}^\infty(\Omega)$ is the complex vector space of all infinitely differentiable functions $\Omega \to \mathbb{C}$. This is a subspace of $\mathcal{C}^k(\Omega)$ for all k.

Definition 19.2.8. Let $n \in \mathbb{N}$. Let Ω be an open subset of $\mathbb{R}^n$. Then $\mathcal{P}(\Omega)$ is the complex vector space of all complex polynomials of n variables, considered as functions $\Omega \to \mathbb{C}$. This is a subspace of $\mathcal{C}^\infty(\Omega)$.

Proposition 19.2.9. *The space of all convergent sequences in $\mathbb{C}$ is a subspace of the space of all bounded sequences in $\mathbb{C}$, which is a subspace of $\mathbb{C}^\mathbb{N}$.*

PROOF: Easy. $\square$

Definition 19.2.10. Let p be a real number, $p \geq 1$. Let l^p be the set of all complex sequences (z_n) such that $\sum_{n=1}^\infty |z_n|^p < \infty$.

Proposition 19.2.11. *For p a real number ≥ 1, we have that l^p is a subspace of $\mathbb{C}^\mathbb{N}$.*

PROOF:
$\langle 1 \rangle 1$. For all $(x_n), (y_n) \in l^p$, we have $(x_n + y_n) \in l^p$.
  PROOF: From Minkowski's Inequality.
$\langle 1 \rangle 2$. For all $\lambda \in \mathbb{C}$ and $(x_n) \in l^p$ we have $(\lambda x_n) \in l^p$
  PROOF:
$$\sum_{n=1}^\infty |\lambda x_n|^p = |\lambda|^p \sum_{n=1}^\infty |x_n|^p < \infty$$

$\square$

Definition 19.2.12 (Linear Combination)**.** Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a vector space over K. Let $x, x_1, \ldots, x_n \in V$. Then x is a *linear combination* of $x_1, \ldots, x_n$ iff there exist $\alpha_1, \ldots, \alpha_n \in K$ such that

$$x = \alpha_1 x_1 + \cdots + \alpha_n x_n \ .$$

Definition 19.2.13 (Linearly Independent)**.** A finite set of vectors $\{x_1, \ldots, x_n\}$ is *linearly independent* iff, whenever $\alpha_1 x_1 + \cdots + \alpha_n x_n = 0$, then $\alpha_1 = \cdots = \alpha_n = 0$.

A set of vectors is *linearly independent* iff every finite subset is linearly independent; otherwise, it is *linearly dependent.*

Definition 19.2.14 (Span)**.** Let $\mathcal{A}$ be a set of vectors. The *span* of $\mathcal{A}$, $\operatorname{span} \mathcal{A}$, is the set of all linear combinations of elements of $\mathcal{A}$.

Proposition 19.2.15. $\operatorname{span} \mathcal{A}$ *is the smallest subspace of V that includes $\mathcal{A}$.*

PROOF: Easy. □

Definition 19.2.16 (Basis)**.** A *basis* for V is a linearly independent set of vectors $\mathcal{B}$ such that $\operatorname{span} \mathcal{B} = V$.

Definition 19.2.17 (Finite Dimensional)**.** A vector space is *finite dimensional* iff it has a finite basis; otherwise it is *infinite dimensional.*

Proposition 19.2.18. *In a finite dimensional vector space, any two bases have the same number of elements.*

Definition 19.2.19 (Dimension)**.** The *dimension* of a finite dimensional vector space V, $\dim V$, is the number of elements in any basis.

Proposition 19.2.20.
$$\dim K^n = n$$

PROOF: The standard basis is the set of vectors with one coordinate 1 and all others 0. □

Proposition 19.2.21. *The dimension of $\mathbb{C}^n$ as a real vector space is $2n$.*

19.3   Linear Transformations

Definition 19.3.1 (Linear Transformation)**.** Let U and V be vector spaces over K. Let $T : U \to V$. Then T is a *linear transformation* iff

$$\forall \alpha, \beta \in K. \forall x, y \in U. T(\alpha x + \beta y) = \alpha T(x) + \beta T(y) \ .$$

Proposition 19.3.2. *If $T : U \to V$ is a linear transformation then $T(U)$ is a subspace of V.*

Proposition 19.3.3. *If $T : U \to V$ is a linear transformation then the graph of T, $\{(x, T(x)) \mid x \in U\}$, is a subspace of $U \times V$.*

Definition 19.3.4 (Null Space)**.** Let U and V be vector spaces over K and $T : U \to V$. The *null space* of T is

$$\mathcal{N}(T) := \{x \in U \mid T(x) = 0\} \ .$$

Proposition 19.3.5. *If $T : U \to V$ is a linear transformation then $\mathcal{N}(T)$ is a subspace of U.*

Proposition 19.3.6. *Let U and V be vector spaces over K. The set of all linear transformations $U \to V$ is a vector space over K under*

$$(S + T)(u) = S(u) + T(u)$$
$$(\lambda S)(u) = \lambda S(u)$$

19.4   Normed Spaces

Definition 19.4.1 (Norm)**.** Let K be either $\mathbb{R}$ or $\mathbb{C}$. A *norm* on a vector space V over K is a function $\| \ \| : V \to \mathbb{R}$ such that:

1. $\forall x \in V. \|x\| = 0 \Rightarrow x = 0$

2. $\forall \lambda \in K. \forall x \in V. \|\lambda x\| = |\lambda| \|x\|$

3. *Triangle Inequality* $\forall x, y \in V. \|x + y\| \leq \|x\| + \|y\|$

Proposition 19.4.2. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a normed space over K. Define $d : V^2 \to \mathbb{R}$ by $d(x, y) = \|x - y\|$. Then d is a metric on V.*

PROOF: Easy. □

We identify any normed space V with this metric space.

Proposition 19.4.3. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a vector space over K. Let $\| \ \|$ be a norm on V. Then*

$$\|0\| = 0 \ .$$

PROOF:
$$\begin{aligned}
\|0\| &= \|0 \cdot 0\| && \text{(Proposition 19.1.4)} \\
&= |0| \|0\| && \text{(Axiom 2 for a norm)} \\
&= 0 && \qquad\qquad □
\end{aligned}$$

Proposition 19.4.4. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a vector space over K. Let $\| \ \|$ be a norm on V. Let $x \in V$. Then*

$$\|x\| \geq 0 \ .$$

PROOF:

$$
\begin{aligned}
0 &= \|0\| && \text{(Proposition 19.4.3)} \\
&= \|x - x\| \\
&\leq \|x\| + \|-x\| && \text{(Triangle Inequality)} \\
&= \|x\| + \|x\| && \text{(Axiom 2 for a norm)} \\
&= 2\|x\| && \square
\end{aligned}
$$

Proposition 19.4.5. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a vector space over K. Let $\|\ \|$ be a norm on V. Let $x, y \in V$. Then*

$$| \|x\| - \|y\| | \leq \|x - y\| \ .$$

PROOF:
$\langle 1\rangle 1.\ \|x\| - \|y\| \leq \|x - y\|$
    PROOF: $\|x\| \leq \|x - y\| + \|y\|$ by the Triangle Inequality.
$\langle 1\rangle 2.\ \|y\| - \|x\| \leq \|x - y\|$
    PROOF:

$$
\begin{aligned}
\|x\| + \|x - y\| &= \|x\| + \|y - x\| && \text{(Axiom 2 for a norm)} \\
&\leq \|y\| && \text{(Triangle Inequality)}
\end{aligned}
$$

$\square$

Corollary 19.4.5.1. *Let V be a normed space. Then $\|\ \| : V \to \mathbb{R}$ is continuous.*

Definition 19.4.6 (Euclidean Norm)**.** The *Euclidean norm* on $\mathbb{C}^n$ is defined by

$$\|(z_1, \ldots, z_n)\| = \sqrt{|z_1|^2 + \cdots + |z_n|^2}$$

Proposition 19.4.7. *Define $\|\ \| : \mathbb{C}^n \to \mathbb{R}$ by*

$$\|(z_1, \ldots, z_n)\| = |z_1| + \cdots + |z_n|$$

Then this defines a norm on $\mathbb{C}^n$.

PROOF: Easy. $\square$

Proposition 19.4.8. *Define $\|\ \| : \mathbb{C}^n \to \mathbb{R}$ by*

$$\|(z_1, \ldots, z_n)\| = \max(|z_1|, \ldots, |z_n|)$$

Then this defines a norm on $\mathbb{C}^n$.

PROOF: Easy. $\square$

Proposition 19.4.9. *Let Ω be a closed bounded subset of $\mathbb{R}^n$. Define $\|\ \| : \mathcal{C}(\Omega) \to \mathbb{R}$ by $\|f\| = \max_{x \in \Omega} |f(x)|$. Then $\|\ \|$ defines a norm on $\mathcal{C}(\Omega)$.*

PROOF: Easy. $\square$

Proposition 19.4.10. *Let p be a real number, $p \geq 1$. Define $\| \; \| : l^p \to \mathbb{R}$ by*

$$\|(z_n)\| = \left(\sum_{n=0}^{\infty} |z_n|^p \right)^{1/p} .$$

Then this defines a norm on l^p.

PROOF: Easy. The triangle inequality is Minkowski's Inequality. □

Definition 19.4.11 (Normed Space)**.** Let K be either $\mathbb{R}$ or $\mathbb{C}$. A *normed space* over K consists of a vector space V over K and a norm on V.

We shall write simply:

- K^n for the normed space K^n under the Euclidean norm

- $\mathcal{C}(\Omega)$ for the normed space $\mathcal{C}(\Omega)$ under the norm $\|f\| = \max_{x \in \Omega} |f(x)|$, called the *uniform convergence norm*

- l^p for the normed space l^p under the norm $\|(z_n)\| = \left(\sum_{n=0}^{\infty} |z_n|^p \right)^{1/p}$.

Proposition 19.4.12. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a normed space over K. If $x_n \to l$ as $n \to \infty$ in V and $\lambda_n \to \lambda$ as $n \to \infty$ in K, then $\lambda_n x_n \to \lambda l$ as $n \to \infty$.*

PROOF:
⟨1⟩1. LET: $\epsilon > 0$
⟨1⟩2. LET: $K = |\lambda| + \epsilon/2\|l\|$
⟨1⟩3. PICK N such that, for all $n \geq N$, we have $|\lambda_n - \lambda| < \epsilon/2\|l\|$ and $\|x_n - l\| < \epsilon/(2K)$
⟨1⟩4. For all $n \geq N$ we have $|\lambda_n| < K$
⟨1⟩5. $\|\lambda_n x_n - \lambda l\| < \epsilon$
    PROOF:

$$\begin{aligned}
\|\lambda_n x_n - \lambda l\| &\leq \|\lambda_n x_n - \lambda_n l\| + \|\lambda_n l - \lambda l\| \\
&= |\lambda_n| \|x_n - l\| + |\lambda_n - \lambda| \|l\| \\
&< K \frac{\epsilon}{2K} + \frac{\epsilon}{2\|l\|} \|l\| \\
&= \epsilon
\end{aligned}$$

□

Proposition 19.4.13. *In a normed space, if $x_n \to l$ and $y_n \to m$ then $x_n + y_n \to l + m$*

PROOF:

$$\begin{aligned}
\|(x_n + y_n) - (l + m)\| &\leq \|x_n - l\| + \|y_n - m\| \\
&\to 0
\end{aligned}$$

□

Definition 19.4.14 (Uniform Convergence). Let Ω be a closed bounded set in $\mathbb{R}^n$. Let (f_n) be a sequence in $\mathcal{C}(\Omega)$ and $f \in \mathcal{C}(\Omega)$. Then (f_n) *converges uniformly* to f if and only if, for every $\epsilon > 0$, there exists N such that $\forall x \in \Omega. \forall n \geq N. |f_n(x) - f(x)| < \epsilon$.

Proposition 19.4.15. *(f_n) converges uniformly to f iff (f_n) converges to f under the uniform convergence norm.*

PROOF: Easy. □

Proposition 19.4.16. *There is no norm on $\mathcal{C}([0,1])$ that induces pointwise convergence.*

PROOF:
$\langle 1 \rangle 1$. LET: $\| \ \|$ be any norm on $\mathcal{C}([0,1])$
$\langle 1 \rangle 2$. For $n \in \mathbb{N}$, define $g_n \in \mathcal{C}([0,1])$ by
$$g_n(t) = \begin{cases} 2^n t & \text{if } 0 \leq t \leq 2^{-n} \\ 2 - 2^n t & \text{if } 2^{-n} \leq t \leq 2^{1-n} \\ 0 & \text{otherwise} \end{cases}$$
$\langle 1 \rangle 3$. For all $n \in \mathbb{N}$ we have $\|g_n\| \neq 0$
$\langle 1 \rangle 4$. For $n \in \mathbb{N}$,
  LET: $f_n = g_n / \|g_n\|$
$\langle 1 \rangle 5$. For all $n \in \mathbb{N}$, $\|f_n\| = 1$
$\langle 1 \rangle 6$. f_n does not converge to 0
$\langle 1 \rangle 7$. $f_n \to 0$ as $n \to \infty$ pointwise.
□

Definition 19.4.17 (Equivalent Norms). Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a vector spaces over K. Then two norms $\| \ \|_1$ and $\| \ \|_2$ are *equivalent* if and only if, for any sequence (x_n) in V and $l \in V$, we have $x_n \to l$ under $\| \ \|_1$ if and only if $x_n \to l$ under $\| \ \|_2$.

Proposition 19.4.18. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a vector spaces over K. Let $\| \ \|_1$ and $\| \ \|_2$ be norms on V. Then $\| \ \|_1$ and $\| \ \|_2$ are equivalent if and only if there exist positive reals α and β such that, for all $x \in V$,*

$$\alpha \|x\|_1 \leq \|x\|_2 \leq \beta \|x\|_1 \tag{19.1}$$

PROOF:
$\langle 1 \rangle 1$. If $\| \ \|_1$ and $\| \ \|_2$ are equivalent then (19.1) holds.
 $\langle 2 \rangle 1$. ASSUME: $\| \ \|_1$ and $\| \ \|_2$ are equivalent.
 $\langle 2 \rangle 2$. There exists $\alpha > 0$ such that, for all $x \in V$, we have $\alpha \|x\|_1 \leq \|x\|_2$
  $\langle 3 \rangle 1$. ASSUME: for a contradiction $\forall \alpha > 0. \exists x \in V. \alpha \|x\|_1 > \|x\|_2$
  $\langle 3 \rangle 2$. For $n \in \mathbb{Z}^+$, choose $x_n \in V$ such that $1/n \|x_n\|_1 > \|x_n\|_2$
  $\langle 3 \rangle 3$. For $n \in \mathbb{Z}^+$,
   LET:
$$y_n = \frac{1}{\sqrt{n}} \frac{x_n}{\|x_n\|_2}$$

⟨3⟩4. $\|y_n\|_2 \to 0$ as $n \to \infty$
⟨3⟩5. For all $n \in \mathbb{Z}^+$, $\|y_n\|_1 > \sqrt{n}$
⟨3⟩6. $\|y_n\| \not\to 0$ as $n \to \infty$
⟨2⟩3. There exists $\beta > 0$ such that, for all $x \in V$, we have $\|x\|_2 \leq \beta\|x\|_1$
   PROOF: Similar.
⟨1⟩2. If (19.1) holds then $\|\ \|_1$ and $\|\ \|_2$ are equivalent.
   ⟨2⟩1. ASSUME: (19.1) holds.
   ⟨2⟩2. LET: (x_n) be a sequence in V and $l \in V$
   ⟨2⟩3. If $x_n \to l$ under $\|\ \|_1$ then $x_n \to l$ under $\|\ \|_2$.
      ⟨3⟩1. ASSUME: $x_n \to l$ und $\|\ \|_1$.
      ⟨3⟩2. LET: $\epsilon > 0$
      ⟨3⟩3. PICK N such that $\forall n \geq N.\|x_n - l\| < \epsilon/\beta$
      ⟨3⟩4. LET: $n \geq N$
      ⟨3⟩5. $\|x_n - l\|_2 < \epsilon$
         PROOF:
$$\|x_n - l\|_2 \leq \beta\|x_n - l\|_1$$
$$< \epsilon$$
⟨2⟩4. If $x_n \to l$ under $\|\ \|_2$ then $x_n \to l$ under $\|\ \|_1$.
   PROOF: Similar.
☐

Proposition 19.4.19. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a normed space over K. If $x_1, \ldots, x_n \in V$ are linearly independent, then there exists $c > 0$ such that, for all $\alpha_1, \ldots, \alpha_n \in K$,*

$$\|\alpha_1 x_1 + \cdots + \alpha_n x_n\| \geq c(|\alpha_1| + \cdots + |\alpha_n|) \ .$$

PROOF:
⟨1⟩1. LET: $B = \{(\beta_1, \ldots, \beta_n) \in K^n \mid |\beta_1| + \cdots + |\beta_n| = 1\}$
⟨1⟩2. LET: $f : B \to \mathbb{R}$ be the function
$$f(\beta_1, \ldots, \beta_n) = \|\beta_1 x_1 + \cdots + \beta_n x_n\| \ .$$
⟨1⟩3. LET: c be the minimum value in $f(B)$
   PROOF: f is continuous and B is compact.
⟨1⟩4. $c > 0$
   PROOF: We never have $f(\beta_1, \ldots, \beta_n) = 0$ by linear independence.
⟨1⟩5. LET: $\alpha_1, \ldots, \alpha_n \in K$
⟨1⟩6. ASSUME: w.l.o.g. $\alpha_1, \ldots, \alpha_n$ are not all zero.
⟨1⟩7. For $i = 1, \ldots, n$,
   LET: $\beta_i = \alpha_i/(|\alpha_1| + \cdots + |\alpha_n|)$
⟨1⟩8. $(\beta_1, \ldots, \beta_n) \in B$
⟨1⟩9. $f(\beta_1, \ldots, \beta_n) \geq c$
⟨1⟩10. $\|\alpha_1 x_1 + \cdots + \alpha_n x_n\| \geq c(|\alpha_1| + \cdots + |\alpha_n|)$
☐

Theorem 19.4.20. *Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a finite dimensional vector space over K. Then any two norms on V are equivalent.*

PROOF:

$\langle 1 \rangle 1$. PICK a basis $\{e_1, \ldots, e_n\}$ for V

$\langle 1 \rangle 2$. LET: $\| \ \|_0 : V \to \mathbb{R}$ be the function
$$\|\alpha_1 e_1 + \cdots + \alpha_n e_n\|_0 = |\alpha_1| + \cdots + |\alpha_n| \ .$$

$\langle 1 \rangle 3$. $\| \ \|_0$ is a norm.

  $\langle 2 \rangle 1$. $\forall x \in V. \|x\|_0 = 0 \Rightarrow x = 0$

    PROOF: If $|\alpha_1| + \cdots |\alpha_n| = 0$ then $\alpha_1 = \cdots = \alpha_n = 0$.

  $\langle 2 \rangle 2$. $\forall \lambda \in K. \forall x \in V. \|\lambda x\| = |\lambda| \|x\|$

    PROOF:
$$\begin{aligned}
\|\lambda(\alpha_1 e_1 + \cdots + \alpha_n e_n)\|_0 &= \|\lambda\alpha_1 e_1 + \cdots + \lambda\alpha_n e_n\|_0 \\
&= |\lambda\alpha_1| + \cdots + |\lambda\alpha_n| \\
&= |\lambda|(|\alpha_1| + \cdots + |\alpha_n|) \\
&= |\lambda| \|\alpha_1 e_1 + \cdots + \alpha_n e_n\|
\end{aligned}$$

  $\langle 2 \rangle 3$. The triangle inequality holds.

    PROOF:
$$\begin{aligned}
\|(\alpha_1 e_1 + \cdots + \alpha_n e_n) + (\beta_1 e_1 + \cdots + \beta_n e_n)\| &= \|(\alpha_1 + \beta_1)e_1 + \cdots + (\alpha_n + \beta_n)e_n\| \\
&= |\alpha_1 + \beta_1| + \cdots + |\alpha_n + \beta_n| \\
&\leq (|\alpha_1| + \cdots + |\alpha_n|) + (|\beta_1| + \cdots + |\beta_n|) \\
&= \|\alpha_1 e_1 + \cdots + \alpha_n e_n\|_0 + \|\beta_1 e_1 + \cdots + \beta_n e_n\|_0
\end{aligned}$$

$\langle 1 \rangle 4$. LET: $\| \ \|$ be any norm on V.

    PROVE: $\| \ \|$ is equivalent to $\| \ \|_0$.

$\langle 1 \rangle 5$. For all $\alpha_1, \ldots, \alpha_n \in K$,
$$\|\alpha_1 e_1 + \cdots + \alpha_n e_n\| \leq \max(\|e_1\|, \ldots, \|e_n\|)(|\alpha_1| + \cdots |\alpha_n|)$$

  $\langle 2 \rangle 1$. LET: $\alpha_1, \ldots, \alpha_n \in K$

  $\langle 2 \rangle 2$. $\|\alpha_1 e_1 + \cdots + \alpha_n e_n\| \leq \max(\|e_1\|, \ldots, \|e_n\|)(|\alpha_1| + \cdots |\alpha_n|)$

    PROOF:
$$\|\alpha_1 e_1 + \cdots + \alpha_n e_n\| \leq |\alpha_1| \|e_1\| + \cdots + |\alpha_n| \|e_n\|$$
$$\leq (|\alpha_1| + \cdots + |\alpha_n|) \max(\|e_1\|, \ldots, \|e_n\|) \qquad \square$$

$\langle 1 \rangle 6$. LET: $\beta = \max(\|e_1\|, \ldots, \|e_n\|)$

$\langle 1 \rangle 7$. For all $x \in V$,
$$\|x\| \leq \beta \|x\|_0 \ .$$

$\langle 1 \rangle 8$. There exists $\alpha > 0$ such that, for all $x \in V$,
$$\alpha \|x\|_0 \leq \|x\| \ .$$

  PROOF: Proposition 19.4.19.

$\square$

Definition 19.4.21 (Closed Ball)**.** Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a vector space over K. Let $x \in V$ and let r be a positive real number. The *closed ball* with *centre* x and *radius* r is

$$\overline{B(x,r)} := \{y \in V \mid \|x - y\| \leq r\} \ .$$

Definition 19.4.22 (Sphere)**.** Let K be either $\mathbb{R}$ or $\mathbb{C}$. Let V be a vector space over K. Let $x \in V$ and let r be a positive real number. The *sphere* with *centre* x and *radius* r is

$$S(x,r) := \{y \in V \mid \|x - y\| = r\} \ .$$

Proposition 19.4.23. *Every closed ball is closed.*

Proposition 19.4.24. *Every sphere is closed.*

Proposition 19.4.25. *The union of two closed sets is closed.*

Proposition 19.4.26. *The intersection of a nonempty set of closed sets is closed.*

Proposition 19.4.27. *In a normed space V, both $\emptyset$ and V are closed.*

Proposition 19.4.28. *Let V be a normed space and $C \subseteq V$. Then C is closed iff, for every sequence (x_n) in C and $l \in V$, if $x_n \to l$ then $l \in C$.*

PROOF:
$\langle 1 \rangle 1$. If C is closed then, for every sequence (x_n) in C and $l \in V$, if $x_n \to l$ then $l \in C$.
   $\langle 2 \rangle 1$. ASSUME: C is closed.
   $\langle 2 \rangle 2$. LET: (x_n) be a sequence in C.
   $\langle 2 \rangle 3$. LET: $l \in V$
   $\langle 2 \rangle 4$. ASSUME: $x_n \to l$ as $n \to \infty$
   $\langle 2 \rangle 5$. ASSUME: for a contradiction $l \notin C$
   $\langle 2 \rangle 6$. PICK $\epsilon > 0$ such that $B(l, \epsilon) \subseteq V - C$
   $\langle 2 \rangle 7$. PICK N such that $\forall n \geq N . \|x_n - l\| < \epsilon$
   $\langle 2 \rangle 8$. $x_N \in C$
   $\langle 2 \rangle 9$. $\|x_N - l\| < \epsilon$
   $\langle 2 \rangle 10$. Q.E.D.
     PROOF: This is a contradiction.
$\langle 1 \rangle 2$. If, for every sequence (x_n) in C and $l \in V$, if $x_n \to l$ then $l \in C$, then C is closed.
   $\langle 2 \rangle 1$. ASSUME: for every sequence (x_n) in C and $l \in V$, if $x_n \to l$ then $l \in C$.
   $\langle 2 \rangle 2$. LET: $x \in V - C$
   $\langle 2 \rangle 3$. ASSUME: for a contradiction there is no $\epsilon > 0$ such that $B(x, \epsilon) \subseteq V - C$
   $\langle 2 \rangle 4$. For $n \in \mathbb{Z}^+$, PICK $x_n \in B(x, \epsilon) \cap C$
   $\langle 2 \rangle 5$. $x_n \to x$ as $n \to \infty$
   $\langle 2 \rangle 6$. Q.E.D.
     PROOF: This is a contradiction.
◻

Definition 19.4.29 (Closure). Let V be a normed space and $A \subseteq V$. The *closure* of A, $\operatorname{cl} A$, is the intersection of all the closed sets that include A.

Theorem 19.4.30. *Let V be a normed space and $S \subseteq V$. Then the closure of S is the set of all limits of convergent sequences in S.*

PROOF:
$\langle 1 \rangle 1$. For all $l \in \operatorname{cl} S$, there exists a sequence (x_n) that converges to l.
   $\langle 2 \rangle 1$. LET: $l \in \operatorname{cl} S$
   $\langle 2 \rangle 2$. For $n \in \mathbb{N}$, PICK $x_n \in B(l, 1/(n+1)) \cap S$

$\langle 3 \rangle$1. Assume: for a contradiction $B(l, 1/(n+1))$ does not intersect S.

$\langle 3 \rangle$2. $V - B(l, 1/(n+1))$ is a closed set that includes S.

$\langle 3 \rangle$3. cl $S \subseteq V - B(l, 1/(n+1))$

$\langle 3 \rangle$4. $l \notin B(l, 1/(n+1))$

$\langle 3 \rangle$5. Q.E.D.

  Proof: This is a contradiction.

$\langle 2 \rangle$3. $x_n \to l$

$\langle 1 \rangle$2. For every sequence (x_n) in S, if $x_n \to l$ then $l \in$ cl S.

  Proof: Proposition 19.4.28.

☐

Definition 19.4.31 (Dense). Let V be a normed space and $S \subseteq V$. Then S is *dense* iff cl $S = V$.

Proposition 19.4.32. *In $\mathcal{C}([a,b])$, the set of polynomials is dense.*

Proof: By the Weierstrass Theorem. ☐

Proposition 19.4.33. *For any real $p \geq 1$, the set of all sequences with only finitely many non-zero terms is dense in l^p.*

Proof:

$\langle 1 \rangle$1. Let: $p \geq 1$

$\langle 1 \rangle$2. Let: $(z_n) \in l^p$

$\langle 1 \rangle$3. Let: $\epsilon > 0$

$\langle 1 \rangle$4. Pick N such that $\forall n \geq N.|z_n| < \epsilon/2$

$\langle 1 \rangle$5. Let: (y_n) be the sequence with $y_n = z_n$ for $n < N$, and $y_n = 0$ for $n \geq N$

$\langle 1 \rangle$6. $\|(z_n) - (y_n)\| \leq \epsilon/2$

$\langle 1 \rangle$7. $\|(z_n) - (y_n)\| < \epsilon$

☐

Theorem 19.4.34. *Let V be a normed space. Let $S \subseteq V$. Then the following are equivalent:*

  1. S is dense.

  2. For all $x \in V$, there exists a sequence (x_n) in S such that $x_n \to x$.

  3. Every nonempty open subset of V intersects S.

Proof:

$\langle 1 \rangle$1. $1 \Leftrightarrow 2$

  Proof: Theorem 19.4.30.

$\langle 1 \rangle$2. $1 \Leftrightarrow 3$

  Proof:

      S is dense $\Leftrightarrow$ the only closed set that includes S is V

          $\Leftrightarrow$ the only open set that does not intersect S is empty

☐

Definition 19.4.35 (Compact)**.** Let V be a normed space and $S \subseteq V$. Then S is *compact* if and only if every sequence in S has a subsequence that converges to a limit in S.

Proposition 19.4.36. *Every compact set is closed.*

PROOF:
$\langle 1 \rangle 1$. LET: V be a normed space.
$\langle 1 \rangle 2$. LET: $C \subseteq V$ be compact.
$\langle 1 \rangle 3$. LET: (x_n) be a sequence in C that converges to $l \in V$.
$\langle 1 \rangle 4$. PICK a subsequence (y_n) of (x_n) that converges to $m \in C$.
$\langle 1 \rangle 5$. $l = m$
$\langle 1 \rangle 6$. $l \in C$
$\langle 1 \rangle 7$. Q.E.D.
   PROOF: Proposition 19.4.28.
☐

Definition 19.4.37 (Bounded)**.** Let V be a normed space and $S \subseteq V$. Then S is *bounded* iff there exists $r > 0$ such that $S \subseteq B(0, r)$.

Proposition 19.4.38. *In $\mathbb{R}^n$ and $\mathbb{C}^n$, the compact sets are the closed bounded sets.*

PROOF:
$\langle 1 \rangle 1$. LET: $C \subseteq K^n$
$\langle 1 \rangle 2$. If C is compact then C is closed.
   PROOF: Proposition 19.4.36.
$\langle 1 \rangle 3$. If C is compact then C is bounded.
   $\langle 2 \rangle 1$. ASSUME: C is compact.
   $\langle 2 \rangle 2$. ASSUME: for a contradiction C is not bounded.
   $\langle 2 \rangle 3$. For $n \in \mathbb{N}$, PICK $x_n \in C$ with $\|x_n\| > n + 1$.
   $\langle 2 \rangle 4$. PICK a convergent subsequence (x_{n_r}) that converges to $l \in C$
   $\langle 2 \rangle 5$. $\|x_{n_r}\| \to \|l\|$
   $\langle 2 \rangle 6$. $\|x_{n_r}\| \to +\infty$
   $\langle 2 \rangle 7$. Q.E.D.
     PROOF: This is a contradiction.
$\langle 1 \rangle 4$. If C is closed and bounded then C is compact.
   PROOF: By the Bolzano-Weierstrass Theorem.
☐

Proposition 19.4.39. *Let V be a normed space and $S \subseteq V$. Then S is bounded if and only if, for every sequence (x_n) in S and every sequence (λ_n) in K, if $\lambda_n \to 0$ then $\|\lambda_n x_n\| \to 0$.*

PROOF:
$\langle 1 \rangle 1$. If S is bounded then, for every sequence (x_n) in S and every sequence (λ_n) in K, if $\lambda_n \to 0$ then $\|\lambda_n x_n\| \to 0$.
   $\langle 2 \rangle 1$. ASSUME: S is bounded.

$\langle 2 \rangle 2$. PICK $r > 0$ such that $S \subseteq B(0, r)$.
$\langle 2 \rangle 3$. LET: (x_n) be a sequence in S.
$\langle 2 \rangle 4$. LET: (λ_n) be a sequence in K.
$\langle 2 \rangle 5$. ASSUME: $\lambda_n \to 0$
$\langle 2 \rangle 6$. LET: $\epsilon > 0$
$\langle 2 \rangle 7$. PICK N such that $\forall n \geq N.|\lambda_n| < \epsilon/r$
$\langle 2 \rangle 8$. $\forall n \geq N.\|\lambda_n x_n\| < \epsilon$
$\langle 1 \rangle 2$. If S is unbounded then there exists a sequence (x_n) in S and (λ_n) in K
        such that $\lambda_n \to 0$ and $\|\lambda_n x_n\| \not\to 0$.
$\langle 2 \rangle 1$. S is unbounded.
$\langle 2 \rangle 2$. For $n \in \mathbb{N}$, PICK $x_n \in S$ such that $\|x_n\| > n$.
$\langle 2 \rangle 3$. For $n \in \mathbb{N}$,
        LET: $\lambda_n = 1/n$ if $n > 0$, 1 if $n = 0$
$\langle 2 \rangle 4$. $\lambda_n \to 0$
$\langle 2 \rangle 5$. $\|\lambda_n x_n\| > 1$ for all $n > 1$

☐

Proposition 19.4.40. *In $\mathcal{C}([0, 1])$, the unit ball $\overline{B(0, 1)}$ is closed and bounded but not compact.*

PROOF:
$\langle 1 \rangle 1$. $\overline{B(0, 1)}$ is closed.
    PROOF: Proposition 19.4.23.
$\langle 1 \rangle 2$. $\overline{B(0, 1)}$ is bounded.
    PROOF: $\overline{B(0, 1)} \subseteq B(0, 2)$.
$\langle 1 \rangle 3$. $\overline{B(0, 1)}$ is not compact.
    $\langle 2 \rangle 1$. For $n \in \mathbb{N}$,
            LET: $x_n : [0, 1] \to \mathbb{R}$ be the function $x_n(t) = t^n$.
    $\langle 2 \rangle 2$. For $n \in \mathbb{N}$, we have $x_n \in \overline{B(0, 1)}$.
    $\langle 2 \rangle 3$. No subsequence of (x_n) converges.

☐

Theorem 19.4.41 (Riesz's Lemma)**.** *Let X be a closed proper subspace of a normed space V. For every $\epsilon \in (0, 1)$, there exists $x_\epsilon \in V$ such that $\|x_\epsilon\| = 1$ and $\forall x \in X.\|x_\epsilon - x\| \geq \epsilon$.*

PROOF:
$\langle 1 \rangle 1$. PICK $z \in E - X$
$\langle 1 \rangle 2$. LET: $d = \inf_{x \in X} \|z - x\|$
$\langle 1 \rangle 3$. $d > 0$
    $\langle 2 \rangle 1$. PICK- $\epsilon > 0$ such that $B(z, \epsilon) \subseteq E - X$
    $\langle 2 \rangle 2$. $d \geq \epsilon$
$\langle 1 \rangle 4$. For all $\epsilon \in (0, 1)$, choose $y_\epsilon \in X$ such that
$$d \leq \|z - y_\epsilon\| < d/\epsilon \ .$$
$\langle 1 \rangle 5$. For $\epsilon \in (0, 1)$,
        LET:
$$x_\epsilon = \frac{z - y_\epsilon}{\|z - y_\epsilon\|} \ .$$

$\langle 1 \rangle 6$. For all $x \in X$ we have $\|x_\epsilon - x\| > \epsilon$

PROOF:

$$\|x_\epsilon - x\| = \|\frac{z - y_\epsilon}{\|z - y_\epsilon\|} - x\|$$

$$= \frac{1}{\|z - y_\epsilon\|}\|z - y_\epsilon - \|z - y_\epsilon\|x\| \qquad (y_\epsilon + \|z - y_\epsilon\|x \in X)$$

$$\geq \frac{1}{\|z - y_\epsilon\|}d$$

$$> \epsilon$$

▯

Theorem 19.4.42. *Let V be a normed space. Then V is finite dimensional if and only if $\overline{B(0,1)}$ is compact.*

PROOF:

$\langle 1 \rangle 1$. If V is finite dimensional then $\overline{B(0,1)}$ is compact.

$\quad \langle 2 \rangle 1$. ASSUME: V is finite dimensional.

$\quad \langle 2 \rangle 2$. PICK a basis $\{e_1, \ldots, e_n\}$.

$\quad \langle 2 \rangle 3$. ASSUME: w.l.o.g. $\forall \alpha_1, \ldots, \alpha_n \in K. \|\alpha_1 e_1 + \cdots + \alpha_n e_n\| = |\alpha_1| + \cdots + |\alpha_n|$

$\quad \langle 2 \rangle 4$. LET: (x_m) be a sequence in $\overline{B(0,1)}$

$\quad \langle 2 \rangle 5$. For $m \in \mathbb{N}$,

$\quad\quad$ LET: $x_m = \alpha_{m1}e_1 + \cdots + \alpha_{mn}e_n$

$\quad \langle 2 \rangle 6$. For $m \in \mathbb{N}$ and $i = 1, \ldots, n$, we have $|\alpha_{mi}| \leq 1$

$\quad \langle 2 \rangle 7$. For $i = 1, \ldots, n$, PICK a convergent subsequence $(\alpha_{m_r i})$ of (α_{mi}) in $\mathbb{C}$ that converges to l_i

$\quad\quad$ PROOF: Since $\overline{B(0,1)}$ is compact in K.

$\quad \langle 2 \rangle 8$. x_m converges to $l_1 e_1 + \cdots + l_n e_n$

$\quad\quad$ PROOF:

$$\|x_m - l_1 e_1 - \cdots - l_n e_n\| = \|(\alpha_{m1} - l_1)e_1 + \cdots + (\alpha_{mn} - l_n)e_n\|$$

$$= |\alpha_{m1} - l_1| + \cdots + |\alpha_{mn} - l_n|$$

$$\to 0 \qquad\qquad \text{as } m \to \infty$$

$\langle 1 \rangle 2$. If V is infinite dimensional then $\overline{B(0,1)}$ is not compact.

$\quad \langle 2 \rangle 1$. ASSUME: V is infinite dimensional.

$\quad \langle 2 \rangle 2$. Choose a sequence (x_n) such that $\|x_n\| = 1$ and $\|x_m - x_n\| \geq 1/2$ for all $m \neq n$.

$\quad\quad \langle 3 \rangle 1$. ASSUME: as induction hypothesis there exists a sequence $(x_0, x_1, \ldots, x_n)$ such that $\|x_i\| = 1$ and $\|x_i - x_j\| \geq 1/2$ for $i \neq j$

$\quad\quad \langle 3 \rangle 2$. PICK x_{n+1} such that $\|x_{n+1}\| = 1$ and $\|x_{n_1} - x\| \geq 1/2$ for $x \in \{x_1, \ldots, x_n\}$.

$\quad \langle 2 \rangle 3$. ASSUME: for a contradiction (x_{n_r}) is a subsequence that converges to l

$\quad \langle 2 \rangle 4$. For all r we have $1/2 \leq \|x_{n_r} - l\| + \|x_{n_{r+1}} - l\|$

$\quad\quad$ PROOF:

$$1/2 \leq \|x_{n_r} - x_{n_{r+1}}\| \qquad\qquad\qquad (\langle 2 \rangle 2)$$

$$\leq \|x_{n_r} - l\| + \|x_{n_{r+1}} - l\| \qquad \text{(Triangle Inequality)}$$

$\quad \langle 2 \rangle 5$. Q.E.D.

PROOF: This is a contradiction.
□

Proposition 19.4.43. *Let U and V be normed spaces. Let $f : U \to V$. If f is continuous at one point, then it is continuous.*

PROOF:
$\langle 1 \rangle 1$. ASSUME: f is continuous at $x_0 \in U$.
$\langle 1 \rangle 2$. LET: (x_n) be a sequence in U that converges to $l \in U$.
$\langle 1 \rangle 3$. $x_n - l + x_0 \to x_0$ as $n \to \infty$
$\langle 1 \rangle 4$. $f(x_n - l + x_0) \to f(x_0)$ as $n \to \infty$
$\langle 1 \rangle 5$. $f(x_n) - f(l) + f(x_0) \to f(x_0)$ as $n \to \infty$
$\langle 1 \rangle 6$. $f(x_n) \to f(l)$ as $n \to \infty$.
□

Definition 19.4.44 (Bounded). Let U and V be normed spaces and $L : U \to V$ be a linear transformation. Then L is *bounded* iff there exists $\alpha > 0$ such that $\forall x \in U.\|L(x)\| \leq \alpha\|x\|$.

Theorem 19.4.45. *Let U and V be normed spaces. Let $L : U \to V$ be a linear transformation. Then L is continuous if and only if it is bounded.*

PROOF:
$\langle 1 \rangle 1$. If L is continuous then L is bounded.
  $\langle 2 \rangle 1$. ASSUME: L is not bounded.
  $\langle 2 \rangle 2$. For $n \in \mathbb{N}$, choose $x_n \in U$ such that $\|L(x_n)\| > (n+1)\|x_n\|$
  $\langle 2 \rangle 3$. For $n \in \mathbb{N}$,
       LET: $y_n = x_n/(n+1)\|x_n\|$
  $\langle 2 \rangle 4$. $y_n \to 0$ as $n \to \infty$
  $\langle 2 \rangle 5$. For $n \in \mathbb{N}$, we have $\|L(y_n)\| > 1$
  $\langle 2 \rangle 6$. $L(y_n) \not\to 0 = L(0)$ as $n \to \infty$
  $\langle 2 \rangle 7$. L is not continuous.
$\langle 1 \rangle 2$. If L is bounded then L is continuous.
  $\langle 2 \rangle 1$. LET: $\alpha > 0$ be such that $\forall x \in U.\|L(x)\| \leq \alpha\|x\|$
       PROVE: L is continuous at 0.
  $\langle 2 \rangle 2$. LET: (x_n) be a sequence in U that converges to 0.
  $\langle 2 \rangle 3$. $L(x_n) \to 0$ as $n \to \infty$
  $\langle 2 \rangle 4$. Q.E.D.
     PROOF: Proposition 19.4.43.
  □
□

Corollary 19.4.45.1. *If U and V are finite dimensional normed spaces, then every linear transformation $U \to V$ is continuous.*

Definition 19.4.46. For U and V normed spaces, let $\mathcal{B}(U, V)$ be the space of all bounded linear transformations $U \to V$. This is a subspace of the space of all linear transformations $U \to V$.

Define a norm on $\mathcal{B}(U,V)$ by

$$\|L\| = \sup\{\|L(x)\| \mid x \in U, \|x\| = 1\} \ .$$

We prove this is a norm.

PROOF:
⟨1⟩1. $\forall L \in \mathcal{B}(U,V).\|L\| = 0 \Rightarrow L = 0$
  ⟨2⟩1. LET: $L \in \mathcal{B}(U,V)$
  ⟨2⟩2. ASSUME: $\|L$
          $= 0$
  ⟨2⟩3. For all $x \in U$, if $\|x\| = 1$ then $\|L(x)\| = 0$
  ⟨2⟩4. LET: $x \in U$
     PROVE: $L(x) = 0$
  ⟨2⟩5. ASSUME: w.l.o.g. $x \neq 0$
  ⟨2⟩6. $\|L(x/\|x\|)\| = 0$
    PROOF: ⟨2⟩3
  ⟨2⟩7. $L(x/\|x\|) = 0$
  ⟨2⟩8. $L(x)/\|x\| = 0$
  ⟨2⟩9. $L(x) = 0$
⟨1⟩2. $\forall \lambda \in K.\forall L \in \mathcal{B}(U,V).\|\lambda L\| = |\lambda|\|L\|$
  ⟨2⟩1. LET: $\lambda \in K$
  ⟨2⟩2. LET: $L \in \mathcal{B}(U,V)$
  ⟨2⟩3. $\|\lambda L\| = |\lambda|\|L\|$
    PROOF:

$$
\begin{aligned}
\|\lambda L\| &= \sup_{\|x\|=1} \|\lambda L(x)\| \\
&= \sup_{\|x\|=1} (|\lambda|\|L(x)\|) \\
&= |\lambda| \sup_{\|x\|=1} \|L(x)\| \\
&= |\lambda|\|L\|
\end{aligned}
$$

⟨1⟩3. The triangle inequality holds.
  ⟨2⟩1. LET: $L, M \in \mathcal{B}(U,V)$
  ⟨2⟩2. For all $x \in U$, if $\|x\| = 1$ then
$$\|L(x) + M(x)\| \leq \|L(x)\| + \|M(x)\| \ .$$
  ⟨2⟩3. For all $x \in U$, if $\|x\| = 1$ then
$$\|L(x) + M(x)\| \leq \|L\| + \|M\| \ .$$
  ⟨2⟩4. $\|L + M\| \leq \|L\| + \|M\|$
□

19.5 Banach Spaces

Definition 19.5.1 (Cauchy Sequence). A sequence (x_n) in a normed space is a *Cauchy sequence* iff, for every $\epsilon > 0$, there exists N such that $\forall m, n \geq N.\|x_m - x_n\| < \epsilon$.

Theorem 19.5.2. *Let V be a normed space. Let (x_n) be a sequence in V. Then the following are equivalent.*

1. *(x_n) is Cauchy.*

2. *For every pair of strictly increasing sequences of natural numbers (p_n) and (q_n), we have $\|x_{p_n} - x_{q_n}\| \to 0$ as $n \to \infty$.*

3. *For every strictly increasing sequence of natural numbers (p_n), we have $\|x_{p_{n+1}} - x_{p_n}\| \to 0$ as $n \to \infty$.*

PROOF:

$\langle 1 \rangle 1$. $1 \Rightarrow 2$

  $\langle 2 \rangle 1$. ASSUME: (x_n) is Cauchy.

  $\langle 2 \rangle 2$. LET: (p_n) and (q_n) be a pair of increasing sequences of natural numbers.

  $\langle 2 \rangle 3$. LET: $\epsilon > 0$

  $\langle 2 \rangle 4$. PICK N such that $\forall m, n \geq N.\|x_m - x_n\| < \epsilon$

  $\langle 2 \rangle 5$. $\forall n \geq N.p_n, q_n \geq N$

  $\langle 2 \rangle 6$. $\forall n \geq N.\|x_{p_n} - x_{q_n}\| < \epsilon$

$\langle 1 \rangle 2$. $2 \Rightarrow 3$

  PROOF: Trivial.

$\langle 1 \rangle 3$. $2 \Rightarrow 1$

  $\langle 2 \rangle 1$. ASSUME: (x_n) is not Cauchy.

  $\langle 2 \rangle 2$. PICK $\epsilon > 0$ such that, for all N, there exist $m, n \geq N$ such that $\|x_m - x_n\| \geq \epsilon$

  $\langle 2 \rangle 3$. PICK strictly increasing sequences of natural numbers (p_n) and (q_n) such that, for all n, $\|x_{p_n} - x_{q_n}\| \geq \epsilon$

    $\langle 3 \rangle 1$. ASSUME: as induction hypothesis we have chosen $(p_1, \ldots, p_n)$ and $(q_1, \ldots, q_n)$ strictly increasing such that $\forall i.\|x_{p_i} - x_{q_i}\| \geq \epsilon$

    $\langle 3 \rangle 2$. PICK $p_{n+1}, q_{n+1} \geq \max(p_n, q_n)$ such that $\|x_{p_{n+1}} - x_{q_{n+1}}\| \geq \epsilon$

  $\langle 2 \rangle 4$. 2 is false.

$\langle 1 \rangle 4$. $3 \Rightarrow 2$

  $\langle 2 \rangle 1$. ASSUME: (p_n) and (q_n) are strictly increasonig sequences such that $\|x_{p_n} - x_{q_n}\| \not\to 0$ as $n \to \infty$.

  $\langle 2 \rangle 2$. PICK $\epsilon > 0$ such that, for all N, there exists $n \geq N$ such that $\|x_{p_n} - x_{q_n}\| \geq \epsilon$

  $\langle 2 \rangle 3$. Choose a strictly increasing sequence (r_n) such that, for all n, we have $\|x_{r_{2n}} - x_{r_{2n+1}}\| \geq \epsilon$

    $\langle 3 \rangle 1$. ASSUME: as induction hypothesis we have chosen $(r_0, r_1, \ldots, r_{2n+1})$ such that, for $i = 0, 1, \ldots, n$, we have $\|x_{r_{2i}} - x_{r_{2i+1}}\| \geq \epsilon$

    $\langle 3 \rangle 2$. PICK $i, j \geq r_{2n+1}$ such that $\|x_i - x_j\| \geq \epsilon$

    $\langle 3 \rangle 3$. Set $r_{2n+2} = \min(i, j)$ and $r_{2n+3} = \max(i, j)$

  $\langle 2 \rangle 4$. 3 is false.

$\square$

Proposition 19.5.3. *Every convergent sequence is a Cauchy sequence.*

PROOF:

$\langle 1 \rangle$1. LET: (x_n) be a convergent sequence in a normed space V with limit l.
$\langle 1 \rangle$2. LET: $\epsilon > 0$
$\langle 1 \rangle$3. PICK N such that $\forall n \geq N.\|x_n - l\| < \epsilon/2$
$\langle 1 \rangle$4. $\forall m, n \geq N.\|x_m - x_n\| < \epsilon$
$\square$

Proposition 19.5.4. *Let $\mathcal{P}([0,1])$ be the space of polynomials on $[0,1]$ under the norm of uniform convergence. For $n \in \mathbb{N}$, let $P_n = 1 + x + x^2/2! + \cdots + x^n/n!$. Then (P_n) is Cauchy but does not converge, since e^x is not a polynomial.*

PROOF: Easy. $\square$

Proposition 19.5.5. *If (x_n) is a Cauchy sequence in a normed space V, then $(\|x_n\|)$ converges in $\mathbb{R}$.*

PROOF:
$\langle 1 \rangle$1. $(\|x_n\|)$ is Cauchy.
    $\langle 2 \rangle$1. LET: $\epsilon > 0$
    $\langle 2 \rangle$2. PICK N such that $\forall m, n \geq N.\|x_m - x_n\| < \epsilon$
    $\langle 2 \rangle$3. $\forall m, n \geq N.|\|x_m\| - \|x_n\|| < \epsilon$
      PROOF: Proposition 19.4.5.
$\langle 1 \rangle$2. Q.E.D.
    PROOF: Since $\mathbb{R}$ is complete.
$\square$

Corollary 19.5.5.1. *Every Cauchy sequence is bounded.*

Definition 19.5.6 (Banach space)**.** A normed space is *complete* or a *Banach space* iff every Cauchy sequence converges.

Proposition 19.5.7. *For all $p \geq 1$, the space l^p in complete.*

PROOF:
$\langle 1 \rangle$1. LET: (a_n) be a Cauchy sequence in l^p.
$\langle 1 \rangle$2. For $n \in \mathbb{N}$,
      LET: $a_n = (\alpha_{n0}, \alpha_{n1}, \ldots)$
$\langle 1 \rangle$3. For all $\epsilon > 0$, there exists N such that $\forall m, n \geq N$

$$\sum_{k=0}^{\infty} |\alpha_{mk} - \alpha_{nk}|^p < \epsilon \ .$$

    PROOF: $\langle 1 \rangle$1
$\langle 1 \rangle$4. For all $k \in \mathbb{N}$ and $\epsilon > 0$, there exists N such that $\forall m, n \geq N$
$$|\alpha_{mk} - \alpha_{nk}| < \epsilon \ .$$
    PROOF: $\langle 1 \rangle$3
$\langle 1 \rangle$5. For all $k \in \mathbb{N}$, the sequence $(\alpha_{nk})_n$ converges in $\mathbb{C}$.
    PROOF: $\mathbb{C}$ is complete.
$\langle 1 \rangle$6. For $k \in \mathbb{N}$,
      LET:
$$\alpha_k = \lim_{n \to \infty} \alpha_{nk} \ .$$

$\langle 1 \rangle 7$. LET: $a = (\alpha_k)_k$

$\langle 1 \rangle 8$. For all $\epsilon > 0$, there exists N such that $\forall n \geq N$

$$\sum_{k=0}^{\infty} |\alpha_k - \alpha_{nk}|^p < \epsilon \ .$$

PROOF: Take the limit $m \to \infty$ in $\langle 1 \rangle 4$.

$\langle 1 \rangle 9$. $a \in l^p$

 $\langle 2 \rangle 1$. PICK N such that $\forall n \geq N.\ \sum_{k=0}^{\infty} |\alpha_k - \alpha_{nk}|^p < 1$

 $\langle 2 \rangle 2$. $a - a_N \in l^p$

 $\langle 2 \rangle 3$. Q.E.D.

 PROOF: Since l^p is closed under $+$.

$\langle 1 \rangle 10$. $a_n \to a$ as $n \to \infty$.

 PROOF: Immediate from $\langle 1 \rangle 8$.

$\Box$

Proposition 19.5.8. *For any real number a, b with $a < b$, the space $\mathcal{C}([a, b])$ is complete.*

PROOF:

$\langle 1 \rangle 1$. LET: (f_n) be a Cauchy sequence in $\mathcal{C}([a, b])$

$\langle 1 \rangle 2$. For all $\epsilon > 0$, there exists N such that, for all $m, n \geq N$ and $x \in [a, b]$,

$$|f_m(x) - f_n(x)| < \epsilon \ .$$

$\langle 1 \rangle 3$. For all $x \in [a, b]$, $(f_n(x))_n$ is Cauchy.

$\langle 1 \rangle 4$. LET: $f : [a, b] \to \mathbb{C}$ be the function

$$f(x) = \lim_{n \to \infty} f_n(x) \ .$$

$\langle 1 \rangle 5$. For all $\epsilon > 0$, there exists N such that, for all $n \geq N$ and $x \in [a, b]$,

$$|f_n(x) - f(x)| < \epsilon \ .$$

PROOF: Take the limit $m \to \infty$ in $\langle 1 \rangle 2$.

$\langle 1 \rangle 6$. f is continuous.

 $\langle 2 \rangle 1$. LET: $x_0 \in [a, b]$

 $\langle 2 \rangle 2$. LET: $\epsilon > 0$

 $\langle 2 \rangle 3$. PICK N such that, for all $n \geq N$ and $y \in [a, b]$, we have

$$|f_n(y) - f(y)| < \epsilon/3 \ .$$

 PROOF: $\langle 1 \rangle 5$

 $\langle 2 \rangle 4$. PICK $\delta > 0$ such that, for all $y \in [a, b]$ with $|x_0 - y| < \delta$, we have $|f_N(x_0) - f_N(y)| < \epsilon/3$

 $\langle 2 \rangle 5$. For all $y \in [a, b]$, if $|x_0 - y| < \delta$ then $|f(x_0) - f(y)| < \epsilon$

 PROOF:

$$|f(x_0) - f(y)|$$

$$\leq |f(x_0) - f_N(x_0)| + |f_N(x_0) - f_N(y)| + |f_N(y) - f(y)| \quad \text{(Triangle Inequality)}$$

$$< \epsilon/3 + \epsilon/3 + \epsilon/3 \quad\quad\quad\quad\quad\quad\quad\quad\quad\quad\quad\quad (\langle 2 \rangle 3, \langle 2 \rangle 4)$$

$$= \epsilon$$

$\langle 1 \rangle 7$. $f_n \to f$ as $n \to \infty$.

 PROOF: Immediate from $\langle 1 \rangle 5$.

$\Box$

Definition 19.5.9 (Convergent Series). Let (x_n) be a sequence in a normed space V. We say the series $\sum_{n=0}^{\infty} x_n$ is *convergent* iff the sequence $(\sum_{n=0}^{N} x_n)_N$ is convergent, and then we write $\sum_{n=0}^{\infty} x_n = l$ for $\lim_{N \to \infty} \sum_{n=0}^{N} x_n = l$.

Definition 19.5.10 (Absolutely Convergent Series). Let (x_n) be a sequence in a normed space V. We say the series $\sum_{n=0}^{\infty} x_n$ is *absolutely convergent* iff the series $\sum_{n=0}^{\infty} \|x_n\|$ converges in $\mathbb{R}$.

Theorem 19.5.11. *A normed space is complete if and only if every absolutely convergent series is convergent.*

PROOF:
$\langle 1 \rangle 1$. LET: V be a normed space.
$\langle 1 \rangle 2$. If V is complete then every absolutely convergent series is convergent.
 $\langle 2 \rangle 1$. ASSUME: V is complete.
 $\langle 2 \rangle 2$. LET: (x_n) be a sequence in V.
 $\langle 2 \rangle 3$. ASSUME: $\sum_{n=0}^{\infty} \|x_n\|$ converges.
   PROVE: $(\sum_{n=0}^{N} x_n)_N$ is Cauchy.
 $\langle 2 \rangle 4$. LET: $\epsilon > 0$
 $\langle 2 \rangle 5$. PICK N such that $\sum_{n=N}^{\infty} \|x_n\| < \epsilon$
 $\langle 2 \rangle 6$. LET: $m > n \geq N$
 $\langle 2 \rangle 7$. $\|\sum_{k=0}^{m} x_k - \sum_{k=0}^{n} x_k\| < \epsilon$
   PROOF:

$$\| \sum_{k=0}^{m} x_k - \sum_{k=0}^{n} x_k \| = \| \sum_{k=n+1}^{m} x_k \|$$
$$\leq \sum_{k=n+1}^{m} \|x_k\|$$
$$\leq \sum_{k=n+1}^{\infty} \|x_k\|$$
$$< \epsilon$$

$\langle 1 \rangle 3$. If every absolutely convergent series is convergent then V is complete.
 $\langle 2 \rangle 1$. ASSUME: Every absolutely convergent series is convergent.
 $\langle 2 \rangle 2$. LET: (x_n) be a Cauchy sequence in V.
 $\langle 2 \rangle 3$. Choose an increasing sequence of natural numbers (p_k) such that, for all $m, n \geq p_k$, we have
$$\|x_m - x_n\| < 2^{-k} \ .$$
 $\langle 2 \rangle 4$. $\sum_{k=0}^{\infty} \|x_{p_{k+1}} - x_{p_k}\|$ is absolutely convergent.
 $\langle 2 \rangle 5$. $\sum_{k=0}^{\infty} \|x_{p_{k+1}} - x_{p_k}\|$ is convergent.
 $\langle 2 \rangle 6$. (x_{p_k}) converges.
 $\langle 2 \rangle 7$. LET: $l = \lim_{k \to \infty} x_{p_k}$
 $\langle 2 \rangle 8$. $x_n \to l$ as $n \to \infty$.
   PROOF:
$$\|x_n - l\| \leq \|x_n - x_{p_n}\| + \|x_{p_n} - l\|$$
$$\to 0 \qquad\qquad \text{as } n \to \infty \text{(Theorem 19.5.2)}$$

☐

Proposition 19.5.12. *A closed subspace of a Banach space is a Banach space.*

PROOF:
⟨1⟩1. LET: V be a Banach space.
⟨1⟩2. LET: U be a closed subspace of V.
⟨1⟩3. LET: (x_n) be a Cauchy sequence in U.
⟨1⟩4. (x_n) is a Cauchy sequence in V.
⟨1⟩5. LET: l be the limit of (x_n) in V.
⟨1⟩6. $l \in U$
    PROOF: Proposition 19.4.28.
☐

Definition 19.5.13 (Completion). Let V be a normed space. A *completion* of V consists of a Banach space W and a function $\phi : V \to W$ such that:

- ϕ is injective.

- ϕ is a linear transformation.

- ϕ preserves the norm.

- $\phi(V)$ is dense in W.

Definition 19.5.14 (Equivalent Cauchy Sequences). Let V be a normed space. Let (x_n) and (y_n) be Cauchy sequences in V. Then (x_n) and (y_n) are *equivalent*, $(x_n) \sim (y_n)$, iff $\|x_n - y_n\| \to 0$ as $n \to \infty$.

Proposition 19.5.15. *Equivalence is an equivalence relation on the set of Cauchy sequences.*

Proposition 19.5.16. *If $(x_n) \sim (y_n)$ then $\lim_{n\to\infty} \|x_n\| = \lim_{n\to\infty} \|y_n\|$.*

Theorem 19.5.17. *Let V be a normed space. Let W be the quotient set of all Cauchy sequences modulo $\sim$. Define $+$, $\cdot$ and $\|\ \|$ on W by*

$$[(x_n)] + [(y_n)] = [(x_n + y_n)]$$
$$\lambda[(x_n)] = [(\lambda x_n)]$$
$$\|[(x_n)]\| = \lim_{n\to\infty} \|x_n\|$$

Define $\phi : V \to W$ by $\phi(x)$ is the constant sequence (x). Then $\phi : V \to W$ is the completion of V.

PROOF:
⟨1⟩1. $+$, $\cdot$ and $\|\ \|$ are well defined.
⟨1⟩2. W is a normed space.
⟨1⟩3. $\phi(V)$ is dense in W.
    PROOF: For any $[(x_n)] \in W$ we have $[(x_n)] = \lim_{n\to\infty} \phi(x_n)$.
⟨1⟩4. W is complete.

⟨2⟩1. LET: (X_n) be a Cauchy sequence in W.

⟨2⟩2. For all $n \in \mathbb{N}$, PICK $x_n \in V$ such that $\|\phi(x_n) - X_n\| < 1/n$

PROOF: ⟨1⟩3

⟨2⟩3. For all m, n we have $\|x_n - x_m\| \le \|X_n - X_m\| + 1/n + 1/m$

PROOF:

$$\begin{aligned}
\|x_n - x_m\| &= \|\phi(x_n) - \phi(x_m)\| \\
&\le \|\phi(x_n) - X_n\| + \|X_n - X_m\| + \|\phi(x_m) - X_m\| \\
&\le \|X_n - X_m\| + 1/n + 1/m
\end{aligned}$$

⟨2⟩4. (x_n) is a Cauchy sequence in V.

⟨2⟩5. LET: $X = [(x_n)]$

PROVE: $X_n \to X$ as $n \to \infty$

⟨2⟩6. $\|X_n - X\| \to 0$ as $n \to \infty$

PROOF:

$$\begin{aligned}
\|X_n - X\| &\le \|X_n - \phi(x_n)\| + \|\phi(x_n) - X\| \\
&< \|\phi(x_n) - X\| + 1/n \\
&\to 0
\end{aligned}$$

⟨1⟩5. ϕ is injective.

⟨1⟩6. ϕ is a linear transformation.

⟨1⟩7. ϕ preserves the norm.

∎