# Mathematics

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## Chapter 1

## Sets and Classes

### 1.1 Classes

Our language is the language of first-order logic with equality over one primitive binary predicate  $\in$ . We call all the objects we reason about *sets*. When  $a \in b$ , we say a is a *member* or *element* of b, or b contains a. We write  $b \ni a$  for  $a \in b$ , and  $a \notin b$  for  $\neg(a \in b)$ . We write  $\forall x \in a.\phi$  as an abbreviation for  $\forall x(x \in a \to \phi)$ , and  $\exists x \in a.\phi$  as an abbreviation for  $\exists x(x \in a \land \phi)$ .

We shall speak informally of *classes* as an abbreviation for talking about predicates. A *class* is determined by a unary predicate  $\phi[x]$  (possibly with parameters). We write  $\{x \mid \phi[x]\}$  or  $\{x : \phi[x]\}$  for the class determined by  $\phi[x]$ . We write 'a is an element of  $\{x \mid \phi[x]\}$ ' or ' $a \in \{x \mid \phi[x]\}$ ' for  $\phi[a]$ .

We write  $\{t[x_1, ..., x_n] \mid P[x_1, ..., x_n]\}$  for

$$\{y \mid \exists x_1, \dots, x_n (y = t[x_1, \dots, x_n] \land P[x_1, \dots, x_n])\}$$
.

We say two classes **A** and **B** are *equal*, and write  $\mathbf{A} = \mathbf{B}$ , iff  $\forall x (x \in \mathbf{A} \leftrightarrow x \in \mathbf{B})$ .

**Proposition Schema 1.1.1.** For any class **A**, the following is a theorem.

$$\mathbf{A} = \mathbf{A}$$

PROOF: We have  $\forall x (x \in \mathbf{A} \Leftrightarrow x \in \mathbf{A})$ .  $\square$ 

**Proposition Schema 1.1.2.** For any classes **A** and **B**, the following is a theorem.

If 
$$\mathbf{A} = \mathbf{B}$$
 then  $\mathbf{B} = \mathbf{A}$ .

PROOF: If  $\forall x (x \in \mathbf{A} \Leftrightarrow x \in \mathbf{B})$  then  $\forall x (x \in \mathbf{B} \Leftrightarrow x \in \mathbf{A})$ .

**Proposition Schema 1.1.3.** For any classes A, B and C, the following is a theorem.

If 
$$A = B$$
 and  $B = C$  then  $A = C$ .

PROOF: If  $\forall x (x \in \mathbf{A} \Leftrightarrow x \in \mathbf{B})$  and  $\forall x (x \in \mathbf{B} \Leftrightarrow x \in \mathbf{C})$  then  $\forall x (x \in \mathbf{A} \Leftrightarrow x \in \mathbf{C})$ .  $\Box$ 

### 1.1.1 Subclasses

**Definition 1.1.4** (Subclass). We say a class **A** is a *subclass* of **B**, or **B** is a *superclass* of **A**, or **B** *includes* **A**, and write  $\mathbf{A} \subseteq \mathbf{B}$  or  $\mathbf{B} \supseteq \mathbf{A}$ , iff every element of **A** is an element of **B**. Otherwise we write  $\mathbf{A} \not\subseteq \mathbf{B}$  or  $\mathbf{B} \not\supseteq \mathbf{A}$ .

We say **A** is a *proper* subclass of **B**, **B** is a *proper* superclass of **A**, or **B** properly includes **A**, and write  $\mathbf{A} \subsetneq \mathbf{B}$  or  $\mathbf{B} \supsetneq \mathbf{A}$ , iff  $\mathbf{A} \subseteq \mathbf{B}$  and  $\mathbf{A} \neq \mathbf{B}$ .

**Proposition Schema 1.1.5.** For any class **A**, the following is a theorem.

$$\mathbf{A} \subseteq \mathbf{A}$$

PROOF: Every element of **A** is an element of **A**.  $\square$ 

**Proposition Schema 1.1.6.** For any classes **A** and **B**, the following is a theorem.

If 
$$A \subseteq B$$
 and  $B \subseteq A$  then  $A = B$ .

PROOF: If every element of A is an element of B, and every element of B is an element of A, then A and B have exactly the same elements.  $\Box$ 

**Proposition Schema 1.1.7.** For any classes A, B and C, the following is a theorem.

If 
$$A \subseteq B$$
 and  $B \subseteq C$  then  $A \subseteq C$ .

PROOF: If every element of A is an element of B and every element of B is an element of C then every element of A is an element of C.

### 1.1.2 Constructions of Classes

**Definition 1.1.8** (Empty Class). The *empty class*  $\emptyset$  is  $\{x \mid \bot\}$ . Every other class is *nonempty*.

**Definition 1.1.9** (Universal Class). The universal class V is  $\{x \mid \top\}$ .

**Definition 1.1.10** (Enumeration). Given objects  $a_1, \ldots, a_n$ , we define the class  $\{a_1, \ldots, a_n\}$  to be the class  $\{x \mid x = a_1 \lor \cdots \lor x = a_n\}$ .

**Definition 1.1.11** (Intersection). For any classes **A** and **B**, the *intersection*  $\mathbf{A} \cap \mathbf{B}$  is  $\{x \mid x \in \mathbf{A} \land x \in \mathbf{B}\}.$ 

**Definition 1.1.12** (Union). For any classes **A** and **B**, the *union*  $\mathbf{A} \cup \mathbf{B}$  is  $\{x \mid x \in \mathbf{A} \lor x \in \mathbf{B}\}.$ 

**Definition 1.1.13** (Relative Complement). Let **A** and **B** be classes. The *relative complement* of **B** in **A** is the class  $\mathbf{A} - \mathbf{B} := \{x \in \mathbf{A} \mid x \notin \mathbf{B}\}.$ 

**Definition 1.1.14** (Symmetric Difference). For any classes **A** and **B**, the *symmetric difference* is the class  $\mathbf{A} + \mathbf{B} := (\mathbf{A} - \mathbf{B}) \cup (\mathbf{B} - \mathbf{A})$ .

**Definition 1.1.15** (Pairwise disjoint). Let **A** be a class. We say the elements of **A** are *pairwise disjoint* iff, for all  $x, y \in \mathbf{A}$ , if  $x \cap y \neq \emptyset$  then x = y.

## 1.2 Sets and the Axiom of Extensionality

**Definition 1.2.1** (Axiom of Extensionality). The *Axiom of Extensionality* is the statement: if two sets have exactly the same members, then they are equal.

$$\forall x, y (\forall z (z \in x \Leftrightarrow z \in y) \Rightarrow x = y)$$
.

When working in a theory with the Axiom of Extensionality, we may identify a set a with the class  $\{x \mid x \in a\}$ . Our use of the symbols  $\in$  and = is consistent. We say a class  $\mathbf{A}$  is a set iff there exists a set a such that  $a = \mathbf{A}$ ; that is,  $\{x \mid \phi[x]\}$  is a set iff  $\exists a \forall x (x \in a \leftrightarrow \phi[x])$ . Otherwise,  $\mathbf{A}$  is a proper class.

**Definition 1.2.2** (Subset). If A is a set and  $A \subseteq \mathbf{B}$ , we say A is a *subset* of **B**.

**Definition 1.2.3** (Union). The *union* of a class **A** is  $\{x \mid \exists X \in \mathbf{A}.x \in X\}$ . We write  $\bigcup_{P(x)} t(x)$  for  $\bigcup \{t(x) \mid P(x)\}$ .

**Definition 1.2.4** (Intersection). The *intersection* of a class **A** is  $\{x \mid \forall X \in \mathbf{A}.x \in X\}$ . We write  $\bigcap_{P(x)} t(x)$  for  $\bigcap \{t(x) \mid P(x)\}$ .

**Definition 1.2.5** (Power Class). For any class **A**, the *power class*  $\mathcal{P}$ **A** is  $\{X \mid X \subseteq \mathbf{A}\}$ .

## 1.3 The Other Axioms

**Definition 1.3.1** (Pairing Axiom). The *Pairing Axiom* is the statement: for any sets a and b, the class  $\{a, b\}$  is a set.

$$\forall a \forall b \exists c \forall x (x \in c \Leftrightarrow x = a \lor x = b)$$

**Definition 1.3.2** (Union Axiom). The *Union Axiom* is the statement: for any set A, the class  $\bigcup A$  is a set.

$$\forall A \exists B \forall x (x \in B \Leftrightarrow \exists y (y \in A \land x \in y))$$

**Definition 1.3.3** (Comprehension Axiom Scheme). The *Comprehension Axiom Scheme* is the set of sentences of the form, for any class A: If A is a subclass of a set then A is a set.

That is, for any property  $P[x, y_1, \ldots, y_n]$ :

For any sets  $a_1, \ldots, a_n$  and B, the class  $\{x \in B \mid P[x, a_1, \ldots, a_n]\}$  is a set.

$$\forall a_1, \dots, a_n, B. \exists C. \forall x (x \in C \Leftrightarrow x \in B \land P[x, a_1, \dots, a_n])$$

**Definition 1.3.4** (Replacement Axiom Scheme). The Replacement Axiom Scheme is the set of sentences of the form, for some property  $P[x, y, z_1, \ldots, z_n]$ :

For any sets  $a_1, \ldots, a_n, B$ , assume for all  $x \in B$  there exists at most one y such that  $P[x, y, a_1, \ldots, a_n]$ . Then  $\{y \mid \exists x \in B.P[x, y, a_1, \ldots, a_n] \text{ is a set. }$ 

$$\forall a_1, \dots, a_n, B(\forall x \in B. \forall y, y'(P[x, y, a_1, \dots, a_n] \land P[x, y', a_1, \dots, a_n] \Rightarrow y = y') \Rightarrow \exists C \forall y (y \in C \Leftrightarrow \exists x \in B. P[x, y, a_1, \dots, a_n]))$$

**Definition 1.3.5** (Power Set Axiom). The *Power Set Axiom* is the statement: the power class of a set is a set.

$$\forall A \exists B \forall x (x \in B \Leftrightarrow \forall y (y \in x \Rightarrow y \in A))$$

**Definition 1.3.6** (Axiom of Infinity). The *Axiom of Infinity* is the statement: there exists a set I such that  $\emptyset \in I$  and  $\forall x \in I.x \cup \{x\} \in I$ .

$$\exists I (\exists e \in I. \forall x. x \notin e \land \forall x \in I. \exists y \in I. \forall z (z \in y \Leftrightarrow z \in x \lor z = x))$$

**Definition 1.3.7** (Axiom of Choice). The *Axiom of Choice* is the statement: For any set A of pairwise disjoint, nonempty sets, there exists a set C such that, for all  $x \in A$ , we have  $x \cap C$  has exactly one element.

$$\forall A(\forall x \in A. \exists yy \in x \land \forall x, y \in A. \forall z(z \in x \land z \in y \Rightarrow x = y) \Rightarrow \exists C. \forall x \in A. \exists y \forall z(z \in x \land z \in C \Leftrightarrow z = y))$$

**Definition 1.3.8** (Axiom of Regularity). The *Axiom of Regularity* is the statement: for any A, if A has a member, then there exists  $m \in A$  such that  $m \cap A = \emptyset$ .

$$\forall A(\exists x. x \in A \Rightarrow \exists m \in A. \neg \exists x (x \in m \land x \in A))$$

**Definition 1.3.9** (Zermelo Set Theory). *Zermelo set theory* is the theory whose axioms are:

- Extensionality
- Pairing
- Union
- Comprehension
- Power Set
- Infinity
- Choice
- Regularity

We label theorems with Z when they are provable in Zermelo set theory.

**Definition 1.3.10** (Zermelo-Fraenkel Set Theory). Zermelo-Fraenkel set theory is the theory whose axioms are:

- Extensionality
- Union

- Replacement
- Power Set
- Infinity
- Choice
- Regularity

We label theorems with ZFC when they are provable in Zermelo-Fraenkel set theory.

We label a theorem with FOL if it can be proved in first-order logic, i.e. from no axioms.

### 1.4 ZFC Extends Z

**Proposition 1.4.1** (Z,ZFC). The empty class  $\emptyset$  is a set.

PROOF: Immediate from the Axiom of Infinity.

Proposition 1.4.2 (ZFC). The Axiom of Pairing is a theorem of ZFC.

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Proof:
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\begin{array}{l} \langle 1 \rangle 1. \text{ Let: } a,b \text{ be sets.} \\ \langle 1 \rangle 2. \text{ Let: } P(x,y) \text{ be the predicate } (x=\emptyset \wedge y=a) \vee (x=\mathcal{P}\emptyset \wedge y=b). \\ \langle 1 \rangle 3. \text{ For all } x \in \mathcal{PP}\emptyset, \text{ there exists at most one } y \text{ such that } P(x,y). \\ \langle 2 \rangle 1. \text{ Let: } x \in \mathcal{PP}\emptyset \\ \langle 2 \rangle 2. \text{ Let: } y \text{ and } y' \text{ be sets.} \\ \langle 2 \rangle 3. \text{ Assume: } P(x,y) \text{ and } P(x,y') \\ \langle 2 \rangle 4. \ (x=\emptyset \wedge y=a) \vee (x=\mathcal{P}\emptyset \wedge y=b) \\ \text{PROOF: From } \langle 2 \rangle 3. \\ \langle 2 \rangle 5. \ (x=\emptyset \wedge y'=a) \vee (x=\mathcal{P}\emptyset \wedge y'=b) \\ \text{PROOF: From } \langle 2 \rangle 3. \\ \langle 2 \rangle 6. \ \emptyset \neq \mathcal{P}\emptyset \\ \text{PROOF: Since } \emptyset \in \mathcal{P}\emptyset \text{ and } \emptyset \notin \emptyset. \\ \langle 2 \rangle 7. \ y=y' \\ \langle 1 \rangle 4. \text{ Let: } A \text{ be the set } \{y \mid \exists x \in \mathcal{PP}\emptyset.P(x,y)\}. \\ \langle 1 \rangle 5. \ A=\{a,b\} \\ \sqcap \end{array}
```

**Proposition Schema 1.4.3** (ZFC). Every instance of the Comprehension Axiom Scheme is a theorem of ZFC.

### Proof:

- $\langle 1 \rangle 1$ . Let: P(x) be a predicate.
- $\langle 1 \rangle 2$ . Let: A be a set.
- $\langle 1 \rangle 3$ . Let: Q(x,y) be the predicate  $P(x) \wedge y = x$ .

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\langle 1 \rangle 4. For all x \in A, there exists at most one y such that Q(x,y).
    \langle 2 \rangle 1. Let: x \in A
    \langle 2 \rangle 2. Let: y and y' be sets.
    \langle 2 \rangle 3. Assume: Q(x,y) and Q(x,y')
    \langle 2 \rangle 4. \ x \in A \land P(x) \land y = x \land y' = x
       PROOF: From \langle 2 \rangle 3.
    \langle 2 \rangle 5. \ y = y'
       PROOF: From \langle 2 \rangle 4.
\langle 1 \rangle 5. Let: B be the set \{ y \mid \exists x \in A.Q(x,y) \}
   PROOF: This is a set by an Axiom of Replacement and \langle 1 \rangle 4.
\langle 1 \rangle 6. \ B = \{ y \in A \mid P(y) \}
   Proof:
                         y \in B \Leftrightarrow \exists x \in A.Q(x,y)
                                                                                                  (\langle 1 \rangle 5)
                                    \Leftrightarrow \exists x \in A(P(x) \land y = x)
                                                                                                  (\langle 1 \rangle 3)
                                    \Leftrightarrow P(y)
```

Corollary Schema 1.4.3.1 (ZFC). Every axiom of Z is a theorem of ZFC.

It follows that every theorem of Z is a theorem of ZFC.

## 1.5 Consequences of the Axioms

**Proposition 1.5.1** (Z). The union of two sets is a set.

PROOF: Because  $A \cup B = \bigcup \{A, B\}$ .  $\square$ 

**Proposition Schema 1.5.2** (Z). For any number n, the following is a theorem: For any sets  $a_1, \ldots, a_n$ , the class  $\{a_1, \ldots, a_n\} = \{x \mid x = a_1 \lor \cdots \lor x = a_n\}$  is a set.

PROOF: The case n=1 follows from Pairing since  $\{a\}=\{a,a\}$ . If we have proved the theorem for n we have  $\{a_1,\ldots,a_n,a_{n+1}\}=\{a_1,\ldots,a_n\}\cup\{a_{n+1}\}$ .  $\square$ 

**Proposition 1.5.3** (Z). No set is a member of itself.

Corollary 1.5.3.1 (Z). The universal class V is a proper class.

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PROOF: If V is a set then  $V \in V$ , contradicting the Proposition.

**Proposition 1.5.4** (Z). There are no sets a and b such that  $a \in b$  and  $b \in a$ .

### Proof:

- $\langle 1 \rangle 1$ . Let: a and b be any sets.
- $\langle 1 \rangle 2$ . Pick  $m \in \{a, b\}$  such that  $m \cap \{a, b\} = \emptyset$
- $\langle 1 \rangle 3$ . Case: m = a

PROOF: Then  $b \notin a$ .

 $\langle 1 \rangle 4$ . Case: m = b

PROOF: Then  $a \notin b$ .

**Proposition 1.5.5** (Z). The intersection of a set and a class is a set.

PROOF: Immediate from Comprehension.

Proposition 1.5.6 (Z). The relative complement of a class in a set is a set.

[Z]

PROOF: Immediate from Comprehension.

Corollary 1.5.6.1 (Z). The symmetric difference of two sets is a set.

**Proposition 1.5.7** (Z). The intersection of a nonempty class is a set.

### Proof:

- $\langle 1 \rangle 1$ . Let: **A** be a nonempty class.
- $\langle 1 \rangle 2$ . Pick  $B \in \mathbf{A}$
- $\langle 1 \rangle 3. \cap \mathbf{A} \subseteq B$
- $\langle 1 \rangle 4$ .  $\bigcap \mathbf{A}$  is a set.

Proof: By Comprehension.

**Proposition Schema 1.5.8** (FOL). For any classes  ${\bf A}$  and  ${\bf B}$ , the following is a theorem:

If 
$$A \subseteq B$$
 then  $\mathcal{P}A \subseteq \mathcal{P}B$ .

PROOF: Every subset of **A** is a subset of **B**.  $\square$ 

**Proposition Schema 1.5.9** (FOL). For any classes **A** and **B**, the following is a theorem:

If 
$$A \subseteq B$$
 then  $\bigcup A \subseteq \bigcup B$ .

PROOF: If  $x \in X \in \mathbf{A}$  then  $x \in X \in \mathbf{B}$ .  $\square$ 

**Proposition Schema 1.5.10** (Z). For any class **A**, the following is a theorem:

$$\mathbf{A} = \bigcup \mathcal{P} \mathbf{A}$$

Proof:

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\begin{array}{l} \langle 1 \rangle 1. \ \mathbf{A} \subseteq \bigcup \mathcal{P} \mathbf{A} \\ \text{Proof: For all } x \in \mathbf{A} \text{ we have } x \in \{x\} \in \mathcal{P} \mathbf{A}. \\ \langle 1 \rangle 2. \ \bigcup \mathcal{P} \mathbf{A} \subseteq \mathbf{A} \\ \langle 2 \rangle 1. \ \text{Let: } x \in \bigcup \mathcal{P} \mathbf{A} \\ \langle 2 \rangle 2. \ \text{Pick } X \in \mathcal{P} \mathbf{A} \text{ such that } x \in X \\ \langle 2 \rangle 3. \ X \subseteq \mathbf{A} \\ \langle 2 \rangle 4. \ x \in \mathbf{A} \\ \end{array}
```

## 1.6 Transitive Classes

**Definition 1.6.1** (Transitive Class). A class **A** is a *transitive class* iff whenever  $x \in y \in \mathbf{A}$  then  $x \in \mathbf{A}$ .

**Proposition Schema 1.6.2** (FOL). For any class **A**, the following is a theorem:

The following are equivalent.

- 1. A is a transitive class.
- 2.  $\bigcup \mathbf{A} \subseteq \mathbf{A}$
- 3. Every element of A is a subset of A.
- 4.  $\mathbf{A} \subseteq \mathcal{P}\mathbf{A}$

Proof: Immediate from definitions.

**Proposition Schema 1.6.3** (FOL). For any class **A**, the following is a theorem:

If **A** is a transitive class then  $\bigcup \mathbf{A}$  is a transitive class.

### Proof:

- $\langle 1 \rangle 1$ . Assume: **A** is a transitive class.
- $\langle 1 \rangle 2$ . Let:  $x \in y \in \bigcup \mathbf{A}$
- $\langle 1 \rangle 3. \ y \in \mathbf{A}$

PROOF: Since  $\bigcup \mathbf{A} \subseteq \mathbf{A}$  by Proposition 1.6.2.

 $\langle 1 \rangle 4. \ x \in \bigcup \mathbf{A}$ 

**Proposition Schema 1.6.4** (Z). For any class A, the following is a theorem: We have A is a transitive class if and only if  $\mathcal{P}A$  is a transitive class.

#### Proof

- $\langle 1 \rangle 1$ . If **A** is a transitive class then  $\mathcal{P}\mathbf{A}$  is a transitive class.
  - $\langle 2 \rangle 1$ . Assume: **A** is a transitive class.
  - $\langle 2 \rangle 2$ .  $\mathbf{A} \subseteq \mathcal{P} \mathbf{A}$

Proof: Proposition 1.6.2.

 $\langle 2 \rangle 3$ .  $\mathcal{P} \mathbf{A} \subseteq \mathcal{P} \mathcal{P} \mathbf{A}$ 

Proof: Proposition 1.5.8.  $\langle 2 \rangle 4$ .  $\mathcal{P}\mathbf{A}$  is a transitive class. Proof: Proposition 1.6.2.  $\langle 1 \rangle 2$ . If  $\mathcal{P}\mathbf{A}$  is a transitive class then  $\mathbf{A}$  is a transitive class.  $\langle 2 \rangle 1$ . Assume:  $\mathcal{P}\mathbf{A}$  is a transitive class.  $\langle 2 \rangle 2$ .  $\bigcup \mathcal{P} \mathbf{A} \subseteq \mathcal{P} \mathbf{A}$ Proof: Proposition 1.6.2.  $\langle 2 \rangle 3$ .  $\mathbf{A} \subseteq \mathcal{P} \mathbf{A}$ Proof: Proposition 1.5.10.  $\langle 2 \rangle 4$ . **A** is a transitive class. Proof: Proposition 1.6.2. Proposition Schema 1.6.5 (FOL). For any class A, the following is a theo-If every member of A is a transitive set then  $\bigcup A$  is a transitive class. Proof:  $\langle 1 \rangle 1$ . Assume: Every member of **A** is a transitive set.  $\langle 1 \rangle 2$ . Let:  $x \in y \in \bigcup \mathbf{A}$  $\langle 1 \rangle 3$ . PICK  $A \in \mathbf{A}$  such that  $y \in A$ .  $\langle 1 \rangle 4. \ x \in A$ PROOF: Since A is a transitive set.  $\langle 1 \rangle 5. \ x \in \bigcup \mathbf{A}$ **Proposition Schema 1.6.6** (FOL). For any class **A**, the following is a theo-If every member of **A** is a transitive set then  $\bigcap \mathbf{A}$  is a transitive class. Proof:  $\langle 1 \rangle 1$ . Assume: Every member of **A** is a transitive set.  $\langle 1 \rangle 2$ . Let:  $x \in y \in \bigcap \mathbf{A}$ Prove:  $x \in \bigcap \mathbf{A}$  $\langle 1 \rangle 3$ . Let:  $A \in \mathbf{A}$  $\langle 1 \rangle 4. \ y \in A$  $\langle 1 \rangle 5. \ x \in A$ Proof: Since A is a transitive set.

# Chapter 2

## Relations

## 2.1 Ordered Pairs

**Definition 2.1.1** (Ordered Pair). For any sets a and b, the *ordered pair* (a, b) is defined to be  $\{\{a\}, \{a, b\}\}.$ 

**Theorem 2.1.2** (Z). For any sets a, b, c, d, we have (a,b) = (c,d) if and only if a = c and b = d.

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Proof:
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\langle 1 \rangle 1. If (a, b) = (c, d) then a = c and b = d.
    \langle 2 \rangle 1. Assume: \{\{a\}, \{a, b\}\} = \{\{c\}, \{c, d\}\}
    \langle 2 \rangle 2. \cap \{\{a\}, \{a, b\}\} = \bigcap \{\{c\}, \{c, d\}\}\
    \langle 2 \rangle 3. \{a\} = \{c\}
    \langle 2 \rangle 4. a = c
    \langle 2 \rangle 5. \bigcup \{\{a\}, \{a, b\}\} = \bigcup \{\{c\}, \{c, d\}\}\
    \langle 2 \rangle 6. \ \{a,b\} = \{c,d\}
    \langle 2 \rangle 7. b = c or b = d
    \langle 2 \rangle 8. a = d or b = d
    \langle 2 \rangle 9. If b = c and a = d then b = d
        Proof: By \langle 2 \rangle 4.
    \langle 2 \rangle 10. b=d
        PROOF: From \langle 2 \rangle 7, \langle 2 \rangle 8, \langle 2 \rangle 9.
\langle 1 \rangle 2. If a = c and b = d then (a, b) = (c, d).
    PROOF: First-order logic.
```

**Definition 2.1.3** (Cartesian Product). The *Cartesian product* of classes **A** and **B** is the class  $\mathbf{A} \times \mathbf{B} := \{(x,y) \mid x \in \mathbf{A}, y \in \mathbf{B}\}.$ 

**Proposition 2.1.4** (Z). For any sets A and B, the class  $A \times B$  is a set.

PROOF: It is a subset of  $\mathcal{PP}(A \cup B)$ .  $\square$ 

**Proposition Schema 2.1.5** (Z). For any classes A, B and C, the following is a theorem:

$$\mathbf{A} \times (\mathbf{B} \cup \mathbf{C}) = (\mathbf{A} \times \mathbf{B}) \cup (\mathbf{A} \times \mathbf{C})$$

Proof:

$$(x,y) \in \mathbf{A} \times (\mathbf{B} \cup \mathbf{C}) \Leftrightarrow x \in \mathbf{A} \wedge (y \in \mathbf{B} \vee y \in \mathbf{C})$$
$$\Leftrightarrow (x \in \mathbf{A} \wedge y \in \mathbf{B}) \vee (x \in \mathbf{A} \wedge y \in \mathbf{C})$$
$$\Leftrightarrow (x,y) \in (\mathbf{A} \times \mathbf{B}) \cup (\mathbf{A} \times \mathbf{C})$$

**Proposition Schema 2.1.6** (Z). For any classes  ${\bf A}$  and  ${\bf B}$ , the following is a theorem:

If 
$$\mathbf{A} \times \mathbf{B} = \mathbf{A} \times \mathbf{C}$$
 and  $\mathbf{A}$  is nonempty then  $\mathbf{B} = \mathbf{C}$ .

Proof:

- $\langle 1 \rangle 1$ . Pick  $a \in \mathbf{A}$
- $\langle 1 \rangle 2$ . For all x we have  $x \in \mathbf{B}$  iff  $x \in \mathbf{C}$ .

Proof:

$$x \in \mathbf{B} \Leftrightarrow (a, x) \in \mathbf{A} \times \mathbf{B}$$
  
 $\Leftrightarrow (a, x) \in \mathbf{A} \times \mathbf{C}$   
 $\Leftrightarrow x \in \mathbf{C}$ 

**Proposition Schema 2.1.7** (Z). For any classes **A** and **B**, the following is a theorem:

$$\mathbf{A} \times \bigcup \mathbf{B} = \{(a,b) \mid \exists Y \in \mathbf{B}. (a \in \mathbf{A} \land b \in Y)\}\$$

Proof:

$$(x,y) \in A \times \bigcup \mathbf{B} \Leftrightarrow x \in A \land \exists Y \in \mathbf{B}. y \in Y$$
$$\Leftrightarrow \exists Y \in \mathbf{B}(x \in A \land y \in Y)$$

## 2.2 Relations

**Definition 2.2.1** (Relation). A relation  $\mathbf{R}$  between classes  $\mathbf{A}$  and  $\mathbf{B}$  is a subclass of  $\mathbf{A} \times \mathbf{B}$ .

A (binary) relation on **A** is a relation between **A** and **A**. We write  $x\mathbf{R}y$  for  $(x,y) \in \mathbf{R}$ .

### 2.2.1 Identity Functions

**Definition 2.2.2** (Identity Function). For any class A, the *identity function* or *diagonal relation*  $\operatorname{id}_A$  on A is

$$id_{\mathbf{A}} := \{(x, x) \mid x \in \mathbf{A}\} .$$

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### 2.2.2 Inverses

**Definition 2.2.3** (Inverse). The *inverse* of a relation  $\mathbf{R}$  between  $\mathbf{A}$  and  $\mathbf{B}$  is the relation  $\mathbf{R}^{-1}$  between  $\mathbf{B}$  and  $\mathbf{A}$  defined by

$$b\mathbf{R}^{-1}a \Leftrightarrow a\mathbf{R}b$$
.

**Proposition Schema 2.2.4** (Z). For any classes A, B and R, the following is a theorem:

If **R** is a relation between **A** and **B**, we have  $(\mathbf{R}^{-1})^{-1} = \mathbf{R}$ .

Proof:

$$x(\mathbf{R}^{-1})^{-1}y \Leftrightarrow y\mathbf{R}^{-1}x$$
  
 $\Leftrightarrow x\mathbf{R}y$ 

### 2.2.3 Composition

**Definition 2.2.5** (Composition). Let  $\mathbf{R}$  be a relation between  $\mathbf{A}$  and  $\mathbf{B}$ , and  $\mathbf{S}$  be a relation between  $\mathbf{B}$  and  $\mathbf{C}$ . The *composition*  $\mathbf{S} \circ \mathbf{R}$  is the relation between  $\mathbf{A}$  and  $\mathbf{C}$  defined by

$$a(\mathbf{S} \circ \mathbf{R})c \Leftrightarrow \exists b(a\mathbf{R}b \wedge b\mathbf{S}c)$$
.

**Proposition Schema 2.2.6** (Z). For any classes A, B, C, R and S, the following is a theorem:

If  ${\bf R}$  is a relation between  ${\bf A}$  and  ${\bf B}$ , and  ${\bf S}$  is a relation between  ${\bf B}$  and  ${\bf C}$ , then

$$(\mathbf{S} \circ \mathbf{R})^{-1} = \mathbf{R}^{-1} \circ \mathbf{S}^{-1} .$$

Proof:

$$z(\mathbf{S} \circ \mathbf{R})^{-1}x \Leftrightarrow x(\mathbf{S} \circ \mathbf{R})z$$

$$\Leftrightarrow \exists y.(x\mathbf{R}y \wedge y\mathbf{S}z)$$

$$\Leftrightarrow \exists y.(y\mathbf{R}^{-1}x \wedge z\mathbf{S}^{-1}y)$$

$$\Leftrightarrow z(\mathbf{R}^{-1} \circ \mathbf{S}^{-1})x$$

### 2.2.4 Properties of Relaitons

**Definition 2.2.7** (Reflexive). Let **R** be a binary relation on **A**. Then **R** is reflexive on **A** iff  $\forall x \in \mathbf{A}.(x,x) \in \mathbf{R}$ .

**Proposition Schema 2.2.8** (Z). For any classes A and R, the following is a theorem:

If **R** is a reflexive relation on **A** then so is  $\mathbf{R}^{-1}$ .

Proof:

 $\langle 1 \rangle 1$ . Let:  $x \in \mathbf{A}$ 

 $\langle 1 \rangle 2$ .  $x \mathbf{R} x$ 

PROOF: Since  $\mathbf{R}$  is reflexive.

$$\langle 1 \rangle 3. \ x \mathbf{R}^{-1} x$$

**Definition 2.2.9** (Irreflexive). A relation **R** is *irreflexive* iff there is no x such that  $(x, x) \in \mathbf{R}$ .

**Definition 2.2.10** (Symmetric). A relation **R** is *symmetric* iff, whenever  $x\mathbf{R}y$ , then  $y\mathbf{R}x$ .

**Definition 2.2.11** (Antisymmetric). A relation **R** is *antisymmetric* iff, whenever  $x\mathbf{R}y$  and  $y\mathbf{R}x$ , then x=y.

**Proposition Schema 2.2.12** (Z). For any classes A and R, the following is a theorem:

If  $\mathbf{R}$  is an antisymmetric relation on  $\mathbf{A}$  then so is  $\mathbf{R}^{-1}$ .

### Proof:

- $\langle 1 \rangle 1$ . Assume:  $x \mathbf{R}^{-1} y$  and  $y \mathbf{R}^{-1} x$
- $\langle 1 \rangle 2$ .  $y \mathbf{R} x$  and  $x \mathbf{R} y$
- $\langle 1 \rangle 3. \ x = y$

Proof: Since  $\mathbf{R}$  is antisymmetric.

**Definition 2.2.13** (Transitive). A relation **R** is *transitive* iff, whenever  $x\mathbf{R}y$  and  $y\mathbf{R}z$ , then  $x\mathbf{R}z$ .

**Proposition Schema 2.2.14** (Z). For any classes A, B and R, the following is a theorem:

If **R** is a transitive relation between **A** and **B** then  $\mathbf{R}^{-1}$  is transitive.

#### PROOF

- $\langle 1 \rangle 1$ . Assume:  $(x,y), (y,z) \in \mathbf{R}^{-1}$
- $\langle 1 \rangle 2. \ (y, x), (z, y) \in \mathbf{R}$
- $\langle 1 \rangle 3. \ (z, x) \in \mathbf{R}$
- $\langle 1 \rangle 4. \ (x,z) \in \mathbf{R}^{-1}$

**Proposition 2.2.15** (Z). For any relation R on a set A, there exists a smallest transitive relation on A that includes R.

PROOF: The relation is  $\bigcap \{S \in \mathcal{P}A^2 \mid R \subseteq S, S \text{ is transitive}\}$ .  $\square$ 

**Definition 2.2.16** (Transitive Closure). For any relation R on a set A, the transitive closure of R is the smallest transitive relation that includes R.

**Definition 2.2.17** (Minimal). Let **R** be a relation on **A**. An element  $m \in \mathbf{A}$  is *minimal* iff there is no  $x \in \mathbf{A}$  such that  $x\mathbf{R}m$ .

**Definition 2.2.18** (Maximal). Let **R** be a relation on **A**. An element  $m \in \mathbf{A}$  is *maximal* iff there is no  $x \in \mathbf{A}$  such that  $m\mathbf{R}x$ .

## 2.3 n-ary Relations

**Definition Schema 2.3.1.** For any sets  $a_1, \ldots, a_n$ , define the *ordered n-tuple*  $(a_1, \ldots, a_n)$  by

$$(a_1) := a_1$$
  
 $(a_1, \dots, a_n, a_{n+1}) = ((a_1, \dots, a_n), a_{n+1})$ 

**Definition Schema 2.3.2.** An n-ary relation on A is a class of ordered n-tuples all of whose components are in A.

## 2.4 Well Founded Relations

**Definition 2.4.1** (Well Founded). A relation  ${\bf R}$  on a class  ${\bf A}$  is well founded iff:

- for all  $a \in A$ , the class  $\{x \in \mathbf{A} \mid x\mathbf{R}a\}$  is a set;
- every nonempty subset of A has an R-minimal element.

**Proposition 2.4.2** (Z). For any class **A**, the relation  $\{(x,y) \in \mathbf{A}^2 \mid x \in y\}$  is well founded.

Proof:

 $\langle 1 \rangle 1$ . For all  $a \in \mathbf{A}$ , the class  $\{x \in \mathbf{A} \mid x \in a\}$  is a set.

PROOF: It is a subclass of a.

 $\langle 1 \rangle 2$ . Every nonempty subset of **A** has an  $\in$ -minimal element.

 $\langle 2 \rangle 1$ . Let: C be a nonempty subset of **A** 

 $\langle 2 \rangle 2$ . Pick  $m \in C$  such that  $m \cap C = \emptyset$ 

PROOF: Axiom of Regularity.

 $\langle 2 \rangle 3$ . m is  $\in$ -minimal in C.

**Proposition Schema 2.4.3** (Z). For any classes A, B and R, the following is a theorem:

Assume **R** is a well founded relation on **A** and **B**  $\subseteq$  **A** is nonempty. Then **B** has an **R**-minimal element.

Proof:

 $\langle 1 \rangle 1$ . Pick  $b \in \mathbf{B}$ 

 $\langle 1 \rangle 2$ . Let:  $S = \{x \in \mathbf{B} \mid x\mathbf{R}b\}$ 

PROOF: S is a set because it is a subclass of  $\{x \in \mathbf{A} \mid x\mathbf{R}b\}$ .

 $\langle 1 \rangle 3$ . Case:  $S = \emptyset$ 

PROOF: In this case b is an **R**-minimal element of **B**.

 $\langle 1 \rangle 4$ . Case:  $S \neq \emptyset$ 

PROOF: In this cases S has an  $\mathbf{R}$ -minimal element, which is an  $\mathbf{R}$ -minimal element of  $\mathbf{B}$ .

Ш

**Theorem Schema 2.4.4** (Transfinite Induction Principle (Z)). For any classes **A**, **B** and **R**, the following is a theorem:

Assume **R** is a well founded relation on **A** and **B**  $\subseteq$  **A**. Assume that, for all  $t \in$  **A**,

$$\{x \in \mathbf{A} \mid x\mathbf{R}t\} \subseteq \mathbf{B} \Rightarrow t \in \mathbf{B}$$
.

Then  $\mathbf{B} = \mathbf{A}$ .

#### PROOF:

- $\langle 1 \rangle 1$ . Assume: for a contradiction  $\mathbf{B} \neq \mathbf{A}$
- $\langle 1 \rangle 2$ . Pick an **R**-minimal element m of  $\mathbf{A} \mathbf{B}$ .

Proof: Proposition 2.4.3.

 $\langle 1 \rangle 3. \{ x \in \mathbf{A} \mid x\mathbf{R}m \} \subseteq \mathbf{B}$ 

PROOF: By minimality of m.

- $\langle 1 \rangle 4. \ m \in \mathbf{B}$
- $\langle 1 \rangle 5$ . Q.E.D.

Proof: This is a contradiction.

**Theorem 2.4.5** (Z). The transitive closure of a well founded relation on a set is well founded.

### Proof:

- $\langle 1 \rangle 1$ . Let: R be a well founded relation on the set A.
- $\langle 1 \rangle 2$ . Let:  $R^t$  be the transitive closure of R.
- $\langle 1 \rangle 3$ . For any  $x, y \in A$ , if  $xR^ty$  then there exists  $z \in A$  such that zRy. PROOF:  $\{(x,y) \in A^2 \mid \exists z \in A. zRy\}$  is a transitive relation on A that includes R.
- $\langle 1 \rangle 4$ . Let: B be a nonempty subset of A.
- $\langle 1 \rangle$ 5. Pick an R-minimal element b of B.
- $\langle 1 \rangle 6$ . b is  $R^t$ -minimal in B.

PROOF: If there exists x such that  $xR^tb$  then there exists z such that zRb by  $\langle 1 \rangle 3$ .

**Definition 2.4.6** (Initial Segment). Let **R** be a relation on **A** and  $a \in \mathbf{A}$ . The *initial segment* up to a is

$$\operatorname{seg} a := \{ x \in \mathbf{A} \mid x\mathbf{R}a \} .$$

**Theorem Schema 2.4.7** (Transfinite Recursion Theorem Schema (ZFC)). For any classes A, R and any property G[x, y, z], there exists a class F such that, for any class F' the following is a theorem:

Assume that **R** is a well-founded relation on **A**. Assume that, for any f and t, there exists a unique z such that G[f,t,z]. Then  $\mathbf{F}: \mathbf{A} \to \mathbf{V}$  such that, for all  $t \in \mathbf{A}$ , we have  $\mathbf{F} \upharpoonright \operatorname{seg} t$  is a set and

$$G[\mathbf{F} \upharpoonright \operatorname{seg} t, t, \mathbf{F}(t)]$$
 .

If  $\mathbf{F}' : \mathbf{A} \to \mathbf{V}$  satisfies that, for all  $t \in \mathbf{A}$ , we have  $\mathbf{F}' \upharpoonright \operatorname{seg} t$  is a set and  $G[\mathbf{F}' \upharpoonright \operatorname{seg} t, t, \mathbf{F}'(t)]$ , then  $\mathbf{F}' = \mathbf{F}$ .

### Proof:

- $\langle 1 \rangle 1$ . For B a subset of A, let us say a function  $v : B \to V$  is acceptable iff, for all  $x \in B$ , we have  $\operatorname{seg} x \subseteq B$  and  $G[v \upharpoonright \operatorname{seg} x, x, v(x)]$
- $\langle 1 \rangle 2$ . Let: **K** be the class of all acceptable functions.
- $\langle 1 \rangle 3$ . Let:  $\mathbf{F} = \bigcup \mathbf{K}$
- $\langle 1 \rangle 4$ . For all  $B, C \subseteq \mathbf{A}$ , given  $v_1 : B \to \mathbf{V}$  and  $v_2 : C \to \mathbf{V}$  acceptable and  $x \in B \cap C$ , we have  $v_1(x) = v_2(x)$ 
  - $\langle 2 \rangle 1$ . Assume: as transfinite induction hypothesis  $\forall y \mathbf{R} x. y \in B \cap C \Rightarrow v_1(y) = v_2(y)$
  - $\langle 2 \rangle 2$ .  $v_1 \upharpoonright \operatorname{seg} x = v_2 \upharpoonright \operatorname{seg} x$
  - $\langle 2 \rangle 3$ .  $G[v_1 \upharpoonright \operatorname{seg} x, x, v_1(x)]$
  - $\langle 2 \rangle 4$ .  $G[v_2 \upharpoonright \operatorname{seg} x, x, v_2(x)]$
  - $\langle 2 \rangle 5. \ v_1(x) = v_2(x)$
- $\langle 1 \rangle 5$ . **F** is a function.
  - $\langle 2 \rangle 1$ . Assume:  $(x,y),(x,z) \in \mathbf{F}$
  - $\langle 2 \rangle 2$ . PICK acceptable  $v_1: B \to \mathbf{V}$  and  $v_2: C \to \mathbf{V}$  such that  $v_1(x) = y$  and  $v_2(x) = z$
  - $\langle 2 \rangle 3. \ y = z$

Proof: By  $\langle 1 \rangle 4$ .

- $\langle 1 \rangle 6$ . For all  $t \in \text{dom } \mathbf{F}$ , we have  $\mathbf{F} \upharpoonright \text{seg } t$  is a set and  $G[\mathbf{F} \upharpoonright \text{seg } t, t, \mathbf{F}(t)]$ 
  - $\langle 2 \rangle 1$ . Let:  $t \in \text{dom } \mathbf{F}$
  - $\langle 2 \rangle 2$ . Pick an acceptable  $v: A \to \mathbf{V}$  such that  $t \in A$
  - $\langle 2 \rangle 3$ . For all  $y \mathbf{R} x$  we have  $v(y) = \mathbf{F}(y)$
  - $\langle 2 \rangle 4$ . **F**  $\upharpoonright \operatorname{seg} x = v \upharpoonright \operatorname{seg} x$
  - $\langle 2 \rangle 5$ .  $G[v \upharpoonright \operatorname{seg} x, x, v(x)]$
  - $\langle 2 \rangle 6$ .  $G[\mathbf{F} \upharpoonright \operatorname{seg} x, x, \mathbf{F}(x)]$
- $\langle 1 \rangle 7$ . dom  $\mathbf{F} = \mathbf{A}$ 
  - $\langle 2 \rangle 1$ . Let:  $x \in \mathbf{A}$
  - $\langle 2 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall y \mathbf{R} x. y \in \mathbf{A}$
  - $\langle 2 \rangle 3$ . Assume: for a contradiction  $x \notin \text{dom } \mathbf{F}$
  - $\langle 2 \rangle 4$ . **F**  $\upharpoonright \operatorname{seg} x$  is a set

Proof: Axiom of Replacement.

- $\langle 2 \rangle 5$ . **F**  $\upharpoonright \operatorname{seg} x$  is acceptable
- $\langle 2 \rangle 6$ . Let: y be the unique object such that  $G[\mathbf{F} \upharpoonright \operatorname{seg} x, x, y]$
- $\langle 2 \rangle 7$ . **F**  $\upharpoonright \operatorname{seg} x \cup \{(x,y)\}$  is acceptable
- $\langle 2 \rangle 8. \ x \in \text{dom } \mathbf{F}$
- $\langle 2 \rangle 9$ . Q.E.D.

PROOF: This is a contradiction.

- $\langle 1 \rangle 8$ . If  $\mathbf{F}' : \mathbf{A} \to \mathbf{V}$  satisfies the theorem, then  $\mathbf{F}' = \mathbf{F}$ .
  - $\langle 2 \rangle 1$ . Let:  $x \in \mathbf{A}$

PROVE:  $\mathbf{F}'(x) = \mathbf{F}(x)$ 

- $\langle 2 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall y \mathbf{R} x. \mathbf{F}'(y) = \mathbf{F}(y)$
- $\langle 2 \rangle 3. \ \mathbf{F} \upharpoonright x = \mathbf{F}' \upharpoonright x$
- $\langle 2 \rangle 4$ .  $G[\mathbf{F} \upharpoonright x, x, \mathbf{F}(x)]$
- $\langle 2 \rangle 5.$   $G[\mathbf{F}' \upharpoonright x, x, \mathbf{F}'(x)]$
- $\langle 2 \rangle 6. \ \mathbf{F}(x) = \mathbf{F}'(x)$

## Chapter 3

# **Functions**

## 3.1 Functions

**Definition 3.1.1** (Function). A function from **A** to **B** is a relation **F** between **A** and **B** such that, for all  $x \in \mathbf{A}$ , there is only one y such that  $x\mathbf{F}y$ . We denote this y by  $\mathbf{F}(x)$ .

A binary operation on a class **A** is a function  $\mathbf{A}^2 \to \mathbf{A}$ .

**Definition 3.1.2** (Closed). Let  $\mathbf{F} : \mathbf{A} \to \mathbf{A}$  be a function and  $\mathbf{B} \subseteq \mathbf{A}$ . Then  $\mathbf{B}$  is *closed* under  $\mathbf{F}$  iff  $\forall x \in \mathbf{B}.\mathbf{F}(x) \in \mathbf{B}$ .

**Proposition 3.1.3** (Z). For any class **A**, the following is a theorem:

$$\mathrm{id}_A:A\to A$$

PROOF: For all  $x \in \mathbf{A}$ , the only y such that  $(x, y) \in \mathrm{id}_{\mathbf{A}}$  is y = x.  $\square$ 

**Proposition Schema 3.1.4** (Z). For any classes A, B, C, F and G, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{G}: \mathbf{B} \to \mathbf{C}$ . Then  $\mathbf{G} \circ \mathbf{F}: \mathbf{A} \to \mathbf{C}$  and, for all  $x \in \mathbf{A}$ , we have

$$(\mathbf{G} \circ \mathbf{F})(x) = \mathbf{G}(\mathbf{F}(x))$$
.

Proof:

```
 \begin{array}{l} \langle 1 \rangle 1. \  \, \forall x \in \mathbf{A}.(x,\mathbf{G}(\mathbf{F}(x))) \in \mathbf{G} \circ \mathbf{F}) \\ \text{Proof: Because } (x,\mathbf{F}(x)) \in \mathbf{F} \text{ and } (\mathbf{F}(x),\mathbf{G}(\mathbf{F}(x))) \in \mathbf{G}. \\ \langle 1 \rangle 2. \  \, \text{If } (x,z) \in \mathbf{F} \circ \mathbf{G} \text{ then } z = \mathbf{G}(\mathbf{F}(x)) \\ \langle 2 \rangle 1. \  \, \text{Pick } y \in \mathbf{B} \text{ such that } x\mathbf{F}y \text{ and } y\mathbf{G}z \\ \langle 2 \rangle 2. \  \, y = \mathbf{F}(x) \\ \langle 2 \rangle 3. \  \, z = \mathbf{G}(y) \\ \langle 2 \rangle 4. \  \, z = \mathbf{G}(\mathbf{F}(x)) \\ \end{array}
```

**Proposition 3.1.5** (Z). For any set A there exists a function  $F : \mathcal{P}A - \{\emptyset\} \to A$  (a choice function for A) such that, for every nonempty  $B \subseteq A$ , we have  $F(B) \in B$ .

```
Proof:
 \langle 1 \rangle 1. Let: A be a set.
 \langle 1 \rangle 2. Let: \mathcal{A} = \{ \{B\} \times B \mid B \in \mathcal{P}A - \{\emptyset\} \}
 \langle 1 \rangle 3. Every member of \mathcal{A} is nonempty.
 \langle 1 \rangle 4. Any two distinct members of \mathcal{A} are disjoint.
 \langle 1 \rangle5. PICK a set C such that, for all X \in \mathcal{A}, we have C \cap X is a singleton.
    Proof: Axiom of Choice.
 \langle 1 \rangle 6. Let: F = C \cap \bigcup \mathcal{A}
 \langle 1 \rangle 7. \ F : \mathcal{P}A - \{\emptyset\} \to A
     \langle 2 \rangle 1. F is a function.
         (3)1. Let: (B, b), (B, b') \in F
         \langle 3 \rangle 2. \ (B,b), (B,b') \in \{B\} \times B
             PROOF: Since (B, b), (B, b') \in \bigcup A.
         (3)3. (B,b), (B,b') \in C \cap (\{B\} \times B)
         \langle 3 \rangle 4. \ (B,b) = (B,b')
             PROOF: From \langle 1 \rangle 5.
         \langle 3 \rangle 5. \ b = b'
     \langle 2 \rangle 2. dom F = \mathcal{P}A - \{\emptyset\}
        Proof:
                     B \in \operatorname{dom} F
                 \Leftrightarrow \exists b.(B,b) \in F
                 \Leftrightarrow \exists b. ((B,b) \in \bigcup \mathcal{A} \land (B,b) \in C)
                 \Leftrightarrow \exists b. \exists B' \in \mathcal{P}A - \{\emptyset\}. ((B,b) \in \{B'\} \times B' \land (B,b) \in C)
                 \Leftrightarrow B \in \mathcal{P}A - \{\emptyset\} \land \exists b \in B.(B,b) \in C
                 \Leftrightarrow B \in \mathcal{P}A - \{\emptyset\}
                                                                                                                            (\langle 1 \rangle 5)
     \langle 2 \rangle 3. ran F \subseteq A
\langle 1 \rangle 8. For every nonempty B \subseteq A we have F(B) \in B
```

**Proposition 3.1.6** (Z). For any relation R between A and B, there exists a function  $H: A \to B$  such that  $H \subseteq R$  (i.e.  $\forall x \in A.xRH(x)$ ).

```
PROOF: \langle 1 \rangle 1. Let: R be a relation between A and B. \langle 1 \rangle 2. Pick a choice function G for B. \langle 1 \rangle 3. Define H: A \to B by H(x) = G(\{y \mid xRy\}) \langle 1 \rangle 4. H \subseteq R
```

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### 3.1.1 Injective Functions

**Definition 3.1.7** (Injective). A function  $\mathbf{F} : \mathbf{A} \to \mathbf{B}$  is one-to-one, injective or an injection,  $\mathbf{F} : \mathbf{A} \rightarrowtail \mathbf{B}$ , iff, for all  $x, y \in \mathbf{A}$ , if  $\mathbf{F}(x) = \mathbf{F}(y)$ , then x = y.

**Proposition 3.1.8** (Z). For any class A, the following is a theorem:  $id_A : A \to A$  is injective.

PROOF: If  $id_{\mathbf{A}}(x) = id_{\mathbf{A}}(y)$  then immediately x = y.  $\square$ 

**Proposition Schema 3.1.9** (Z). For any classes **A**, **B**, **C**, **F**, **G**, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \rightarrowtail \mathbf{B}$  and  $\mathbf{G}: \mathbf{B} \rightarrowtail \mathbf{C}$ . Then  $\mathbf{G} \circ \mathbf{F}: \mathbf{A} \rightarrowtail \mathbf{C}$ .

### Proof:

- $\langle 1 \rangle 1$ . Let:  $x, y \in \mathbf{A}$
- $\langle 1 \rangle 2$ . Assume:  $(\mathbf{G} \circ \mathbf{F})(x) = (\mathbf{G} \circ \mathbf{F})(y)$
- $\langle 1 \rangle 3. \ \mathbf{G}(\mathbf{F}(x)) = \mathbf{G}(\mathbf{F}(y))$
- $\langle 1 \rangle 4$ .  $\mathbf{F}(x) = \mathbf{F}(y)$

PROOF: Since G is injective.

 $\langle 1 \rangle 5. \ x = y$ 

PROOF: Since  $\mathbf{F}$  is injective.

П

**Proposition 3.1.10** (Z). Let  $F: A \to B$  where A is nonempty. There exists  $G: B \to A$  (a left inverse) such that  $G \circ F = \mathrm{id}_A$  if and only if F is one-to-one.

### Proof:

- $\langle 1 \rangle 1$ . If there exists  $G: B \to A$  such that  $G \circ F = \mathrm{id}_A$  then F is one-to-one.
  - $\langle 2 \rangle 1$ . Assume:  $G: B \to A$  and  $G \circ F = I_A$
  - $\langle 2 \rangle 2$ . Let:  $x, y \in A$
  - $\langle 2 \rangle 3$ . Assume: F(x) = F(y)
  - $\langle 2 \rangle 4. \ x = y$

PROOF: x = G(F(x)) = G(F(y)) = y

- $\langle 1 \rangle 2$ . If F is one-to-one then there exists  $G: B \to A$  such that  $G \circ F = I_A$ .
  - $\langle 2 \rangle 1$ . Assume: F is one-to-one.
  - $\langle 2 \rangle 2$ . Pick  $a \in A$
  - $\langle 2 \rangle 3$ . Let:  $G: B \to A$  be the function defined by: G(b) is the (unique)  $x \in A$  such that F(x) = b if there exists such an x, G(b) = a otherwise.
  - $\langle 2 \rangle 4$ . For all  $x \in A$  we have G(F(x)) = x.

3.1.2 Surjective Functions

**Definition 3.1.11** (Surjective). Let  $F: A \to B$ . We say that F is *surjective*, or maps A onto B, and write  $F: A \twoheadrightarrow B$ , iff for all  $y \in B$  there exists  $x \in A$  such that F(x) = y.

**Proposition Schema 3.1.12** (Z). For any class **A**, the following is a theorem:  $id_{\mathbf{A}} : \mathbf{A} \to \mathbf{A}$  is surjective.

PROOF: For any  $y \in \mathbf{A}$  we have  $\mathrm{id}_{\mathbf{A}}(y) = y$ .  $\square$ 

**Proposition Schema 3.1.13** (Z). For any classes A, B, C, F and G, the following is a theorem:

If  $\mathbf{F} : \mathbf{A} \twoheadrightarrow \mathbf{B}$  and  $\mathbf{G} : \mathbf{B} \twoheadrightarrow \mathbf{C}$ , then  $\mathbf{G} \circ \mathbf{F} : \mathbf{A} \twoheadrightarrow \mathbf{C}$ .

### Proof:

- $\langle 1 \rangle 1$ . Let:  $c \in \mathbf{C}$
- $\langle 1 \rangle 2$ . Pick  $b \in \mathbf{B}$  such that  $\mathbf{G}(b) = c$ .
- $\langle 1 \rangle 3$ . Pick  $a \in \mathbf{A}$  such that  $\mathbf{F}(a) = b$ .
- $\langle 1 \rangle 4. \ (\mathbf{G} \circ \mathbf{F})(a) = c$

**Proposition 3.1.14** (Z). Let  $F: A \to B$ . There exists  $H: B \to A$  (a right inverse) such that  $F \circ H = \operatorname{id}_B$  if and only if F maps A onto B.

### Proof:

- $\langle 1 \rangle 1$ . If F has a right inverse then F is surjective.
  - $\langle 2 \rangle 1$ . Assume: F has a right inverse  $H: B \to A$ .
  - $\langle 2 \rangle 2$ . Let:  $y \in B$
  - $\langle 2 \rangle 3. \ F(H(y)) = y$
  - $\langle 2 \rangle 4$ . There exists  $x \in A$  such that F(x) = y
- $\langle 1 \rangle 2$ . If F is surjective then F has a right inverse.
  - $\langle 2 \rangle 1$ . Assume: F is surjective.
  - $\langle 2 \rangle 2$ . PICK a function  $H: B \to A$  such that  $H \subseteq F^{-1}$  PROOF: Proposition 3.1.6.
  - $\langle 2 \rangle 3. \ F \circ H = \mathrm{id}_B$ 
    - $\langle 3 \rangle 1$ . Let:  $y \in B$
    - $\langle 3 \rangle 2. \ (y, H(y)) \in F^{-1}$
    - $\langle 3 \rangle 3. \ F(H(y)) = y$

### 3.1.3 Bijections

**Definition 3.1.15** (Bijection). Let  $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ . Then  $\mathbf{F}$  is *bijective* or a *bijection*,  $\mathbf{F} : \mathbf{A} \approx \mathbf{B}$ , iff it is injective and surjective.

**Proposition Schema 3.1.16** (Z). For any class A, the following is a theorem: The identity function  $\mathrm{id}_A: A \approx A$  is a bijection.

Proof: Proposition 3.1.8 and 3.1.12.  $\square$ 

**Proposition Schema 3.1.17** (Z). For any classes A, B and F, the following is a theorem:

If  $\mathbf{F} : \mathbf{A} \approx \mathbf{B}$  then  $\mathbf{F}^{-1} : \mathbf{B} \approx \mathbf{A}$ .

### Proof:

- $\langle 1 \rangle 1. \ \mathbf{F}^{-1} : \mathbf{B} \to \mathbf{A}$ 
  - $\langle 2 \rangle 1$ . Let:  $b \in \mathbf{B}$

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 $\langle 2 \rangle 2$ . PICK  $a \in \mathbf{A}$  such that  $\mathbf{F}(a) = b$ .

Proof: Since  $\mathbf{F}$  is surjective.

 $\langle 2 \rangle 3. \ (b,a) \in \mathbf{F}^{-1}$ 

 $\langle 2 \rangle 4$ . If  $(b, a') \in \mathbf{F}^{-1}$  then a' = a.

 $\langle 3 \rangle 1$ . Let:  $a' \in \mathbf{A}$  such that  $(b, a') \in \mathbf{F}^{-1}$ 

 $\langle 3 \rangle 2$ .  $\mathbf{F}(a') = \mathbf{F}(a)$ 

 $\langle 3 \rangle 3. \ a' = a$ 

PROOF: Since **F** is injective.

 $\langle 1 \rangle 2$ .  $\mathbf{F}^{-1}$  is injective.

 $\langle 2 \rangle 1$ . Let:  $x, y \in \mathbf{B}$ 

 $\langle 2 \rangle 2$ . Assume:  $\mathbf{F}^{-1}(x) = \mathbf{F}^{-1}(y)$ 

 $\langle 2 \rangle 3. \ x = y$ 

PROOF:  $x = \mathbf{F}(\mathbf{F}^{-1}(x)) = \mathbf{F}(\mathbf{F}^{-1}(y)) = y$ .

 $\langle 1 \rangle 3$ .  $\mathbf{F}^{-1}$  is surjective.

PROOF: For all  $a \in \mathbf{A}$  we have  $\mathbf{F}^{-1}(\mathbf{F}(a)) = a$ .

П

**Proposition Schema 3.1.18** (Z). For any classes A, B, C, F and G, the following is a theorem:

If  $\mathbf{F} : \mathbf{A} \approx \mathbf{B}$  and  $\mathbf{G} : \mathbf{B} \approx \mathbf{C}$  then  $\mathbf{G} \circ \mathbf{F} : \mathbf{A} \approx \mathbf{C}$ .

Proof: Propositions 3.1.9 and 3.1.13.  $\square$ 

### 3.1.4 Restrictions

**Definition 3.1.19** (Restriction). Let  $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ . Let  $\mathbf{C} \subseteq \mathbf{A}$ . The *restriction* of  $\mathbf{F}$  to  $\mathbf{C}$ , denoted  $\mathbf{F} \upharpoonright \mathbf{C}$ , is the function

$$\mathbf{F} \upharpoonright \mathbf{C} : \mathbf{C} \to \mathbf{B}$$
 
$$(\mathbf{F} \upharpoonright \mathbf{C})(x) = \mathbf{F}(x) \qquad (x \in \mathbf{C})$$

### **3.1.5** Images

**Definition 3.1.20** (Image). Let  $F:A\to B$  and  $C\subseteq A$ . The *image* of C under F is the class

$$\mathbf{F}(\mathbf{C}) := \{ \mathbf{F}(x) \mid x \in \mathbf{C} \} .$$

**Proposition Schema 3.1.21** (Z). For any classes **F**, **A** and **B**, the following is a theorem.

If  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$ , then for any subset  $S \subseteq \mathbf{A}$ , the class  $\mathbf{F}(S)$  is a set.

PROOF: By an Axiom of Replacement.

**Proposition Schema 3.1.22** (Z). For any classes A, B, C and F, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C} \subseteq \mathcal{P}\mathbf{A}$ . Then

$$\mathbf{F}\left(\bigcup\mathbf{C}\right) = \{y \mid \exists X \in \mathbf{C}.y \in \mathbf{F}(X)\}$$

Proof:

$$y \in \mathbf{F}\left(\bigcup \mathbf{C}\right) \Leftrightarrow \exists x \in \bigcup \mathbf{C}. y = \mathbf{F}(x)$$
  
 $\Leftrightarrow \exists x. \exists X. X \in \mathbf{C} \land x \in X \land y = \mathbf{F}(x)$   
 $\Leftrightarrow \exists X \in \mathbf{C}. y \in \mathbf{F}(X)$ 

**Proposition Schema 3.1.23** (Z). For any classes A, B, C, D and F, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C}, \mathbf{D} \subseteq \mathbf{A}$ . Then

$$\mathbf{F}(\mathbf{C} \cup \mathbf{D}) = \mathbf{F}(\mathbf{C}) \cup \mathbf{F}(\mathbf{D})$$
.

Proof:

$$y \in \mathbf{F}(\mathbf{C} \cup \mathbf{D}) \Leftrightarrow \exists x \in \mathbf{C} \cup \mathbf{D}. y = \mathbf{F}(x)$$
  
 $\Leftrightarrow \exists x \in \mathbf{C}. y = \mathbf{F}(x) \lor \exists x \in \mathbf{D}. y = \mathbf{F}(x)$   
 $\Leftrightarrow y \in \mathbf{F}(\mathbf{C}) \cup \mathbf{F}(\mathbf{D})$ 

**Proposition 3.1.24** (Z). For any classes F, A, B, C and D, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C}, \mathbf{D} \subseteq \mathbf{A}$ . Then

$$F(A \cap B) \subseteq F(A) \cap F(B)$$
.

Equality holds if  $\mathbf{F}$  is injective.

Proof:

```
\langle 1 \rangle 1. \mathbf{F}(\mathbf{A} \cap \mathbf{B}) \subseteq \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B})
     \langle 2 \rangle 1. Let: y \in \mathbf{F}(\mathbf{A} \cap \mathbf{B})
     \langle 2 \rangle 2. Pick x \in \mathbf{A} \cap \mathbf{B} such that y = \mathbf{F}(x)
     \langle 2 \rangle 3. \ y \in \mathbf{F}(\mathbf{A})
          PROOF: Since x \in \mathbf{A}.
     \langle 2 \rangle 4. \ y \in \mathbf{F}(\mathbf{B})
          PROOF: Since x \in \mathbf{B}.
\langle 1 \rangle 2. If F is injective then \mathbf{F}(\mathbf{A} \cap \mathbf{B}) = \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B}).
     \langle 2 \rangle 1. Assume: F is injective.
     \langle 2 \rangle 2. Let: y \in \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B})
     \langle 2 \rangle 3. Pick x \in \mathbf{A} such that y = \mathbf{F}(x)
     \langle 2 \rangle 4. Pick x' \in \mathbf{B} such that y = \mathbf{F}(x')
     \langle 2 \rangle 5. \ x = x'
          Proof: \langle 2 \rangle 1
     \langle 2 \rangle 6. \ x \in \mathbf{A} \cap \mathbf{B}
     \langle 2 \rangle 7. \ y \in \mathbf{F}(\mathbf{A} \cap \mathbf{B})
```

**Proposition Schema 3.1.25** (Z). For any classes **F**, **A**, **B**, and **C**, the following is a theorem:

Let  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C} \subseteq \mathcal{P}\mathbf{A}$ . Then

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$$\mathbf{F}\left(\bigcap \mathbf{C}\right) \subseteq \bigcap \{\mathbf{F}(X) \mid X \in \mathbf{A}\}$$
.

Equality holds if **F** is injective and **A** is nonempty.

```
Proof:
```

```
\langle 1 \rangle 1. \ \mathbf{F} (\bigcap \mathbf{A}) \subseteq \bigcap \{ \mathbf{F}(X) \mid X \in \mathbf{A} \}
     \langle 2 \rangle 1. Let: y \in \mathbf{F}(\bigcap \mathbf{A})
     \langle 2 \rangle 2. PICK x \in \bigcap \mathbf{A} such that y = \mathbf{F}(x)
     \langle 2 \rangle 3. Let: X \in \mathbf{A}
     \langle 2 \rangle 4. \ x \in X
     \langle 2 \rangle 5. \ y \in \mathbf{F}(X)
\langle 1 \rangle 2. If F is injective then \mathbf{F}(\bigcap \mathbf{A}) = \bigcap \{ \mathbf{F}(X) \mid X \in \mathbf{A} \}
     \langle 2 \rangle 1. Assume: F is injective.
     \langle 2 \rangle 2. Assume: A is nonempty.
     \langle 2 \rangle 3. Let: y \in \bigcap \{ \mathbf{F}(X) \mid X \in \mathbf{A} \}
     \langle 2 \rangle 4. Pick X_0 \in \mathbf{A}
     \langle 2 \rangle 5. Pick x \in X_0 such that (x,y) \in \mathbf{F}
     \langle 2 \rangle 6. \ x \in \bigcap \mathbf{A}
          \langle 3 \rangle 1. Let: X \in \mathbf{A}
          \langle 3 \rangle 2. PICK x' \in X such that (x', y) \in \mathbf{F}.
          \langle 3 \rangle 3. \ x = x'
               Proof: \langle 2 \rangle 1
           \langle 3 \rangle 4. \ x \in X
     \langle 2 \rangle 7. \ y \in \mathbf{F}(\bigcap \mathbf{A})
```

**Proposition 3.1.26** (Z). For any classes A, B, C, D and F, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C}, \mathbf{D} \subseteq \mathbf{A}$ . Then

$$\mathbf{F}(\mathbf{C}) - \mathbf{F}(\mathbf{D}) \subseteq \mathbf{F}(\mathbf{C} - \mathbf{D}) \ .$$

Equality holds if  $\mathbf{F}$  is injective.

Proof:

$$\begin{array}{l} \langle 1 \rangle 1. \ \mathbf{F(C)} - \mathbf{F(D)} \subseteq \mathbf{F(A-B)} \\ \langle 2 \rangle 1. \ \mathrm{LET:} \ y \in \mathbf{F(A)} - \mathbf{F(B)} \\ \langle 2 \rangle 2. \ \mathrm{PICK} \ x \in \mathbf{A} \ \mathrm{such \ that} \ y = \mathbf{F}(x) \\ \langle 2 \rangle 3. \ x \notin \mathbf{B} \\ \langle 2 \rangle 4. \ x \in \mathbf{A} - \mathbf{B} \\ \langle 2 \rangle 5. \ y \in \mathbf{F(A-B)} \\ \langle 1 \rangle 2. \ \mathrm{If} \ \mathbf{F} \ \mathrm{is} \ \mathrm{injective} \ \mathrm{then} \ \mathbf{F(A)} - \mathbf{F(B)} = \mathbf{F(A-B)} \\ \langle 2 \rangle 1. \ \mathrm{Assume:} \ \mathbf{F} \ \mathrm{is} \ \mathrm{injective}. \\ \langle 2 \rangle 2. \ \mathrm{Let:} \ y \in \mathbf{F(A-B)} \\ \langle 2 \rangle 3. \ \mathrm{PICK} \ x \in \mathbf{A} - \mathbf{B} \ \mathrm{such \ that} \ y = \mathbf{F}(x) \\ \langle 2 \rangle 4. \ y \in \mathbf{F(A)} \\ \langle 2 \rangle 5. \ y \notin \mathbf{F(B)} \end{array}$$

- $\langle 3 \rangle 1$ . Assume: for a contradiction  $y \in \mathbf{F}(\mathbf{B})$
- $\langle 3 \rangle 2$ . Pick  $x' \in \mathbf{B}$  such that  $y = \mathbf{F}(x')$
- $\langle 3 \rangle 3. \ x = x'$

Proof:  $\langle 2 \rangle 1$ 

- $\langle 3 \rangle 4. \ x \in \mathbf{B}$
- $\langle 3 \rangle 5$ . Q.E.D.

PROOF: This contradicts  $\langle 2 \rangle 3$ .

## 3.1.6 Inverse Images

**Definition 3.1.27** (Inverse Image). Let  $F:A\to B$  and  $C\subseteq B$ . Then the *inverse image* of C under F is

$$\mathbf{F}^{-1}(\mathbf{C}) = \{ x \in \mathbf{A} \mid \mathbf{F}(x) \in \mathbf{C} \}$$
.

**Proposition Schema 3.1.28** (Z). For any classes A, B, C and F, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C} \subseteq \mathcal{P}\mathbf{B}$ . Then

$$\mathbf{F}^{-1}\left(\bigcap \mathbf{C}\right) = \bigcap \{\mathbf{F}^{-1}(X) \mid X \in \mathbf{C}\}\ .$$

Proof:

$$x \in \mathbf{F}^{-1}\left(\bigcap \mathbf{C}\right) \Leftrightarrow \mathbf{F}(x) \in \bigcap \mathbf{C}$$
$$\Leftrightarrow \forall X \in \mathbf{C}.\mathbf{F}(x) \in X$$
$$\Leftrightarrow \forall X \in \mathbf{C}.x \in \mathbf{F}^{-1}(X)$$

**Proposition Schema 3.1.29** (Z). For any classes A, B, C, D and F, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C}, \mathbf{D} \subseteq \mathbf{B}$ . Then

$$F^{-1}(C - D) = F^{-1}(C) - F^{-1}(D)$$
.

Proof:

$$x \in \mathbf{F}^{-1}(\mathbf{C} - \mathbf{D}) \Leftrightarrow \mathbf{F}(x) \in \mathbf{C} - \mathbf{D}$$
  
 $\Leftrightarrow \mathbf{F}(x) \in \mathbf{C} \wedge \mathbf{F}(x) \notin \mathbf{D}$   
 $\Leftrightarrow x \in \mathbf{F}^{-1}(\mathbf{C}) \wedge x \in \mathbf{F}^{-1}(\mathbf{D})$   
 $\Leftrightarrow x \in \mathbf{F}^{-1}(\mathbf{C}) - \mathbf{F}^{-1}(\mathbf{D})$ 

### 3.1.7 Function Sets

**Proposition 3.1.30** (ZFC). For any classes  ${\bf B}$  and  ${\bf F}$ , the following is a theorem:

Let A be a set. If  $\mathbf{F}: A \to \mathbf{B}$  then  $\mathbf{F}$  is a set.

PROOF: By an Axiom of Replacement, we have  $R = \{ \mathbf{F}(x) \mid x \in A \}$  is a set. Hence  $\mathbf{F}$  is a set since  $\mathbf{F} \subseteq A \times R$ .  $\square$ 

**Definition 3.1.31** (Dependent Product Class). Let I be a set and let  $\mathbf{H}(i)$  be a class for all  $i \in I$ . We write  $\prod_{i \in I} \mathbf{H}(i)$  for the class of all functions  $f: I \to \bigcup_{i \in I} \mathbf{H}(i)$  such that  $\forall i \in I. f(i) \in \mathbf{H}(i)$ . We write  $\mathbf{B}^I$  for  $\prod_{i \in I} \mathbf{B}$  where  $\mathbf{B}$  does not depend on I.

**Proposition Schema 3.1.32** (ZFC). Let I be a set. Let H(i) be a set for every  $i \in I$ . Then  $\prod_{i \in I} \mathbf{H}(i)$  is a set.

```
PROOF:  \begin{array}{l} \langle 1 \rangle 1. \ \{ \mathbf{H}(i) \mid i \in I \} \text{ is a set.} \\ \text{PROOF: By an Axiom of Replacement.} \\ \langle 1 \rangle 2. \bigcup_{i \in I} \mathbf{H}(i) \text{ is a set.} \\ \langle 1 \rangle 3. \prod_{i \in I} \mathbf{H}(i) \text{ is a set.} \\ \text{PROOF: It is a subset of } \mathcal{P} \left( I \times \bigcup_{i \in I} \mathbf{H}(i) \right). \end{array}
```

**Proposition 3.1.33** (Z). Let I be a set. Let H(i) be a set for all  $i \in I$ . If  $\forall i \in I. H(i) \neq \emptyset$  then  $\prod_{i \in I} H(i) \neq \emptyset$ .

```
Proof:
```

```
\begin{array}{l} \langle 1 \rangle 1. \  \, \text{Assume:} \  \, \forall i \in I.H(i) \neq \emptyset \\ \langle 1 \rangle 2. \  \, \text{Let:} \  \, R = \{(i,x) \mid i \in I, x \in H(i)\} \\ \langle 1 \rangle 3. \  \, \text{Pick a function} \  \, f:I \rightarrow \bigcup_{i \in I} H(i) \  \, \text{such that} \  \, f \subseteq R \\ \text{Proof: Proposition } 3.1.6. \\ \langle 1 \rangle 4. \  \, f \in \prod_{i \in I} H(i) \\ \sqcap \end{array}
```

## 3.2 Equinumerosity

**Definition 3.2.1** (Equinumerous). Sets A and B are equinumerous,  $A \approx B$ , iff there exists a bijection between A and B.

**Theorem 3.2.2** (Cantor 1873 (Z)). No set is equinumerous to its power set.

### Proof:

```
\begin{array}{l} \langle 1 \rangle 1. \text{ Assume: for a contradiction } f:A \approx \mathcal{P}A \\ \langle 1 \rangle 2. \text{ Let: } S = \{x \in A \mid x \notin f(x)\} \\ \langle 1 \rangle 3. \text{ Pick } a \in A \text{ such that } f(a) = S \\ \langle 1 \rangle 4. \ a \in S \text{ if and only if } a \notin S \\ \langle 1 \rangle 5. \text{ Q.E.D.} \\ \text{Proof: This is a contradiction.} \\ & \square \end{array}
```

## 3.3 Domination

**Definition 3.3.1** (Dominate). A set A is dominated by a set B,  $A \leq B$ , iff there exists an injection  $A \rightarrow B$ .

**Proposition 3.3.2** (Z). Given sets A and B, if  $A \neq \emptyset$  or  $B = \emptyset$ , then we have  $A \preceq B$  iff there exists a surjective function  $B \to A$ .

### Proof:

- $\langle 1 \rangle 1$ . If  $A \leq B$  and  $A \neq \emptyset$  then there exists a surjective function  $B \to A$ .
  - $\langle 2 \rangle 1$ . Assume:  $f: A \to B$  be injective.
  - $\langle 2 \rangle 2$ . Pick  $a \in A$
  - $\langle 2 \rangle 3$ . Let:  $g: B \to A$  be the function defined by  $g(b) = f^{-1}(b)$  if  $b \in \operatorname{ran} f$ , and g(b) = a otherwise.
  - $\langle 2 \rangle 4$ . g is surjective.
- $\langle 1 \rangle 2$ . If there exists a surjective function  $B \to A$  then  $A \leq B$ .
  - $\langle 2 \rangle 1$ . Assume: there exists a surjective function  $g: B \to A$
  - $\langle 2 \rangle 2$ .  $\forall a \in A. \exists b \in B. g(b) = a$
  - $\langle 2 \rangle 3$ . Choose a function  $f: A \to B$  such that  $\forall a \in A.g(f(a)) = a$
  - $\langle 2 \rangle 4$ . f is injective.

## Chapter 4

# Equivalence Relations

**Definition 4.0.1** (Equivalence Relation). An *equivalence relation* on a class **A** is a binary relation on **A** that is reflexive, symmetric and transitive.

**Proposition 4.0.2** (Z). Equinumerosity is an equivalence relation on the class of all sets.

PROOF: Propositions 3.1.16, 3.1.17, 3.1.18.

**Definition 4.0.3** (Respects). Let **R** be an equivalence relation on **A** and **F**:  $\mathbf{A} \to \mathbf{B}$ . Then **F** respects **A** iff, whenever  $(x,y) \in \mathbf{R}$ , then  $\mathbf{F}(x) = \mathbf{F}(y)$ .

**Definition 4.0.4** (Equivalence Class). Let  $\mathbf{R}$  be an equivalence relation on  $\mathbf{A}$  and  $a \in \mathbf{A}$ . The *equivalence class* of a modulo  $\mathbf{R}$  is

$$[a]_{\mathbf{R}} := \{x \mid a\mathbf{R}x\} .$$

**Proposition Schema 4.0.5** (Z). For any classes  ${\bf A}$  and  ${\bf R}$ , the following is a theorem.

Assume **R** be an equivalence relation on **A**. Let  $a, b \in \mathbf{A}$ . Then  $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$  if and only if  $a\mathbf{R}b$ .

### Proof:

- $\langle 1 \rangle 1$ . If  $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$  then  $a\mathbf{R}b$ .
  - $\langle 2 \rangle 1$ . Assume:  $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$
  - $\langle 2 \rangle 2$ .  $b\mathbf{R}b$

PROOF: Reflexivity

- $\langle 2 \rangle 3. \ b \in [b]_{\mathbf{R}}$
- $\langle 2 \rangle 4. \ b \in [a]_{\mathbf{R}}$
- $\langle 2 \rangle 5$ .  $a\mathbf{R}b$
- $\langle 1 \rangle 2$ . If  $a\mathbf{R}b$  then  $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$ .
  - $\langle 2 \rangle 1$ . For all  $x, y \in \mathbf{A}$ , if  $x \mathbf{R} y$  then  $[y]_{\mathbf{R}} \subseteq [x]_{\mathbf{R}}$ 
    - $\langle 3 \rangle 1$ . Let:  $x, y \in \mathbf{A}$
    - $\langle 3 \rangle 2$ . Assume:  $x \mathbf{R} y$

```
\langle 3 \rangle 3. \text{ Let: } t \in [y]_{\mathbf{R}}
\langle 3 \rangle 4. y\mathbf{R}t
\langle 3 \rangle 5. x\mathbf{R}t
\text{Proof: Transitivity, } \langle 3 \rangle 2, \langle 3 \rangle 4.
\langle 3 \rangle 6. t \in [x]_{\mathbf{R}}
\langle 2 \rangle 2. \text{ Assume: } a\mathbf{R}b
\langle 2 \rangle 3. [b]_{\mathbf{R}} \subseteq [a]_{\mathbf{R}}
\text{Proof: } \langle 2 \rangle 1, \langle 2 \rangle 2.
\langle 2 \rangle 4. b\mathbf{R}a
\text{Proof: Symmetry, } \langle 2 \rangle 2.
\langle 2 \rangle 5. [a]_{\mathbf{R}} \subseteq [b]_{\mathbf{R}}
\text{Proof: } \langle 2 \rangle 1, \langle 2 \rangle 4.
\langle 2 \rangle 6. [a]_{\mathbf{R}} = [b]_{\mathbf{R}}
\text{Proof: } \langle 2 \rangle 3, \langle 2 \rangle 5.
```

**Definition 4.0.6** (Partition). A partition  $\Pi$  of a set A is a set of nonempty subsets of A that is disjoint and exhaustive, i.e.

- 1. no two different sets in  $\Pi$  have any common elements, and
- 2. each element of A is in some set in  $\Pi$ .

**Definition 4.0.7.** Let R be an equivalence relation on a set A. The *quotient* set A/R is the set of all equivalence classes.

**Theorem 4.0.8** (Z). Let A be a set and  $\mathbf{B}$  a class. Let R be an equivalence relation on A and  $F:A\to \mathbf{B}$ . Then F respects R if and only if there exists  $\hat{F}:A/R\to \mathbf{B}$  such that

$$\forall a \in A. \hat{F}([a]_R) = F(a) .$$

In this case,  $\hat{F}$  is unique.

### Proof:

- $\langle 1 \rangle 1$ . If F respects R then there exists  $\hat{F}: A/R \to \mathbf{B}$  such that  $\forall a \in A.\hat{F}([a]_R) = F(a)$ .
  - $\langle 2 \rangle 1$ . Assume: F respects R.
  - $\langle 2 \rangle 2$ . Let:  $\hat{F} = \{ ([a]_R, F(a)) \mid a \in A \}$
  - $\langle 2 \rangle 3$ .  $\hat{F}$  is a function.
    - $\langle 3 \rangle 1$ . Assume:  $a, a' \in A$  and  $[a]_R = [a']_R$ Prove: F(a) = F(a')
    - $\langle 3 \rangle 2. \ (a, a') \in R$

Proof: Proposition 4.0.5.

 $\langle 3 \rangle 3$ . F(a) = F(a')

Proof:  $\langle 2 \rangle 1$ 

- $\langle 2 \rangle 4$ . dom  $\hat{F} = A/R$
- $\langle 2 \rangle 5$ . ran  $\hat{F} \subseteq \mathbf{B}$

```
\langle 2 \rangle 6. \ \forall a \in A. \hat{F}([a]_R) = F(a)
\langle 1 \rangle 2. If there exists \hat{F}: A/R \to \mathbf{B} such that \forall a \in A.\hat{F}([a]_R) = F(a) then F
        respects R.
   \langle 2 \rangle 1. Assume: \hat{F}: A/R \to \mathbf{B} and \forall a \in A.\hat{F}([a]_R) = F(a)
   \langle 2 \rangle 2. Let: a, a' \in A
   \langle 2 \rangle 3. Assume: (a, a') \in R
   \langle 2 \rangle 4. [a]_R = [a']_R
      Proof: Proposition 4.0.5.
   \langle 2 \rangle 5. F(a) = F(a')
      Proof: \langle 2 \rangle 1
\langle 1 \rangle 3. If G, H : A/R \to \mathbf{B} and \forall a \in A.G([a]_R) = H([a]_R) then G = H.
Proposition 4.0.9 (Z). Let R be an equivalence relation on a set A. Then
A/R is a partition of A.
Proof:
\langle 1 \rangle 1. Every member of A/R is nonempty.
   PROOF: Since a \in [a]_R by reflexivity.
\langle 1 \rangle 2. No two different sets in A/R have any common elements.
   \langle 2 \rangle 1. Let: [a]_R, [b]_R \in A/R
   \langle 2 \rangle 2. Let: c \in [a]_R \cap [b]_R
            PROVE: [a]_R = [b]_R
   \langle 2 \rangle 3. aRc
      Proof: \langle 2 \rangle 2
   \langle 2 \rangle 4. \ bRc
      Proof: \langle 2 \rangle 2
   \langle 2 \rangle 5. cRb
      Proof: Symmetry, \langle 2 \rangle 4
   \langle 2 \rangle 6. aRb
      Proof: Transitivity, \langle 2 \rangle 3, \langle 2 \rangle 5
   \langle 2 \rangle 7. [a]_R = [b]_R
      Proof: Proposition 4.0.5, \langle 2 \rangle 6
\langle 1 \rangle 3. Each element of A is in some set in A/R.
   PROOF: Since a \in [a]_R by reflexivity.
```

**Proposition 4.0.10** (Z). For any partition P of a set A, there exists a unique equivalence relation R on A such that A/R = P, namely xRy iff  $\exists X \in P(x \in X \land y \in X)$ .

Proof: Easy.

**Definition 4.0.11** (Natural Map). Let A be a set and R an equivalence relation on A. The natural map  $A \to A/R$  is the function that maps  $a \in A$  to  $[a]_R$ .

## Chapter 5

# Ordering Relations

## 5.1 Partial Orders

**Definition 5.1.1** (Partial Ordering). Let **A** be a class. A *partial ordering* on **A** is a relation **R** on **A** that is reflexive, antisymmetric and transitive.

We often write  $\leq$  for a partial ordering, and then write x < y for  $x \leq y \land x \neq y$ .

**Proposition Schema 5.1.2** (Z). For any classes A and R, the following is a theorem:

If **R** is a partial order on **A** then so is  $\mathbf{R}^{-1}$ .

```
Proof:
```

```
\begin{array}{c} \langle 1 \rangle 1. \ \mathbf{R}^{-1} \ \text{is reflexive.} \\ \text{Proof: Proposition 2.2.8.} \\ \langle 1 \rangle 2. \ \mathbf{R}^{-1} \ \text{is antisymmetric.} \\ \text{Proof: Proposition 2.2.12.} \\ \langle 1 \rangle 3. \ \mathbf{R}^{-1} \ \text{is transitive.} \\ \langle 2 \rangle 1. \ \text{Assume: } x\mathbf{R}^{-1}y \ \text{and } y\mathbf{R}^{-1}z \\ \langle 2 \rangle 2. \ y\mathbf{R}x \ \text{and } z\mathbf{R}y \\ \langle 2 \rangle 3. \ z\mathbf{R}x \\ \text{Proof: Since } \mathbf{R} \ \text{is transitive.} \\ \langle 2 \rangle 4. \ x\mathbf{R}^{-1}z \\ \square \end{array}
```

**Proposition Schema 5.1.3** (Z). For any classes A, B, F and R, the following is a theorem:

Assume **R** is a partial order on **B** and **F**:  $\mathbf{A} \to \mathbf{B}$  is injective. Define **S** on **A** by  $x\mathbf{S}y$  iff  $\mathbf{F}(x)\mathbf{RF}(y)$ . Then **S** is a partial order on **A**.

### Proof:

 $\langle 1 \rangle 1$ . **S** is reflexive.

PROOF: For any  $x \in \mathbf{A}$  we have  $\mathbf{F}(x)\mathbf{RF}(x)$ .

```
\langle 1 \rangle2. S is antisymmetric.

\langle 2 \rangle1. Let: x, y \in \mathbf{A}

\langle 2 \rangle2. Assume: x\mathbf{S}y and y\mathbf{S}x

\langle 2 \rangle3. \mathbf{F}(x)\mathbf{R}\mathbf{F}(y) and \mathbf{F}(y)\mathbf{R}\mathbf{F}(x)

\langle 2 \rangle4. \mathbf{F}(x) = \mathbf{F}(y)

PROOF: R is antisymmetric.

\langle 2 \rangle5. x = y

\langle 1 \rangle3. S is transitive.
```

**Corollary Schema 5.1.3.1** (Z). For any classes **A**, **B** and **R**, the following is a theorem:

Assume **R** be a partial order on **A** and **B**  $\subseteq$  **A**. Then **R**  $\cap$  **B**<sup>2</sup> is a partial order on **B**.

**Definition 5.1.4** (Partially Ordered Set). A partially ordered set or poset is a pair  $(A, \leq)$  where A is a set and  $\leq$  is a partial ordering on A. We often write just A for  $(A, \leq)$ .

If  $(A, \leq)$  is a poset and  $B \subseteq A$  we write just B for the poset  $(B, \leq \cap B^2)$ .

**Definition 5.1.5** (Strictly Monotone). Let  $(A, <_A)$  and  $(B, <_B)$  be posets. A function  $f: A \to B$  is *strictly monotone* iff, whenever  $x <_A y$ , then  $f(x) <_B f(y)$ .

**Definition 5.1.6** (Least). Let  $\leq$  be a partial order on  $\mathbf{A}$ . An element  $m \in \mathbf{A}$  is *least* iff for all  $x \in \mathbf{A}$  we have  $m \leq x$ .

**Proposition 5.1.7** (Z). A partial order has at most one least element.

PROOF: If m and m' are least then  $m \leq m'$  and  $m' \leq m$ , so m = m'.  $\square$ 

**Definition 5.1.8** (Greatst). Let  $\leq$  be a partial order on **A**. An element  $m \in \mathbf{A}$  is *greatest* iff for all  $x \in A$  we have  $x \leq m$ .

**Proposition 5.1.9** (Z). A poset has at most one greatest element.

PROOF: If m and m' are greatest then  $m \leq m'$  and  $m' \leq m$ , so m = m'.  $\square$ 

**Definition 5.1.10** (Upper Bound). Let  $\leq$  be a partial ordering on **A** and  $\mathbf{B} \subseteq \mathbf{A}$ . Let  $u \in \mathbf{A}$ . Then u is an *upper bound* for **B** iff  $\forall x \in \mathbf{B}.x \leq u$ .

**Definition 5.1.11** (Lower Bound). Let  $\leq$  be a partial ordering on **A** and  $\mathbf{B} \subseteq \mathbf{A}$ . Let  $l \in \mathbf{A}$ . Then l is a *lower bound* for **B** iff  $\forall x \in \mathbf{B}.l \leq x$ .

**Definition 5.1.12** (Bounded Above). Let  $\leq$  be a partial ordering on **A** and  $\mathbf{B} \subseteq \mathbf{A}$ . Then **B** is *bounded above* iff it has an upper bound.

**Definition 5.1.13** (Bounded Below). Let  $\leq$  be a partial ordering on **A** and  $\mathbf{B} \subseteq \mathbf{A}$ . Then **B** is *bounded below* iff it has a lower bound.

**Definition 5.1.14** (Least Upper Bound). Let  $\leq$  be a partial ordering on  $\mathbf{A}$  and  $\mathbf{B} \subseteq \mathbf{A}$ . Let  $s \in \mathbf{A}$ . Then s is the *least upper bound* or *supremum* of  $\mathbf{B}$  iff s is an upper bound for  $\mathbf{B}$  and, for every upper bound u for  $\mathbf{B}$ , we have  $s \leq u$ .

**Definition 5.1.15** (Greatest Lower Bound). Let  $\leq$  be a partial ordering on  $\mathbf{A}$  and  $\mathbf{B} \subseteq \mathbf{A}$ . Let  $i \in \mathbf{A}$ . Then i is the *greatest lower bound* or *infimum* of  $\mathbf{B}$  iff i is a lower bound for  $\mathbf{B}$  and, for every lower bound l for  $\mathbf{B}$ , we have  $i \leq l$ .

**Definition 5.1.16** (Complete). A partial order is *complete* iff every nonempty subset bounded above has a supremum, and every nonempty subset bounded below has an infimum.

**Definition 5.1.17** (Order Isomorphism). Let A and B be posets. An *order isomorphism* between A and B,  $f:A\cong B$ , is a bijection  $f:A\approx B$  such that f and  $f^{-1}$  are monotone.

**Theorem 5.1.18** (Knaster Fixed-Point Theorem (Z)). Let A be a complete poset with a greatest and least element. Let  $\phi: A \to A$  be monotone. Then there exists  $a \in A$  such that  $\phi(a) = a$ .

#### Proof:

```
\langle 1 \rangle 1. Let: B = \{ x \in A \mid x \le \phi(x) \}
\langle 1 \rangle 2. Let: a = \sup B
```

PROOF: B is nonempty because the least element of A is in B, and it is bounded above by the greatest element of A.

```
\langle 1 \rangle3. For all b \in B we have b \le \phi(a)
\langle 2 \rangle1. Let: b \in B
```

$$\langle 2 \rangle 2. \ b \leq \phi(b)$$

$$\langle 2 \rangle 3. \ b \leq a$$

$$\langle 2 \rangle 4. \ \phi(b) \leq \phi(a)$$

$$\langle 2 \rangle 5. \ b \leq \phi(a)$$

$$\langle 1 \rangle 4. \ a \leq \phi(a)$$

$$\langle 1 \rangle 5. \ \phi(a) \le \phi(\phi(a))$$

$$\langle 1 \rangle 6. \ \phi(a) \in B$$

$$\langle 1 \rangle 7. \ \phi(a) \le a$$

$$\langle 1 \rangle 8. \ \phi(a) = a$$

**Definition 5.1.19** (Dense). Let  $\leq$  be a partial order on **A** and **B**  $\subseteq$  **A**. Then **B** is *dense* iff, for all  $x, y \in$  **A**, if x < y then there exists  $z \in$  **B** such that x < z < y.

**Proposition 5.1.20** (Z). Let A be a complete poset with no least element. Let  $B \subseteq A$  be dense. Let  $\theta : A \to A$  be a monotone map that is the identity on B. Then  $\theta = id_A$ .

```
\langle 1 \rangle 1. Let: a \in A
Prove: \theta(a) = a
```

```
\langle 1 \rangle 2. Let: S(a) = \{ b \in B \mid b < a \}
\langle 1 \rangle 3. S(a) is nonempty and bounded above.
   \langle 2 \rangle 1. S(a) is nonempty.
      \langle 3 \rangle 1. Pick a_1 < a
          Proof: Since a is not least.
       \langle 3 \rangle 2. There exists b \in B such that a_1 < b < a.
   \langle 2 \rangle 2. S(a) is bounded above by a.
\langle 1 \rangle 4. sup S(a) \leq a
\langle 1 \rangle 5. sup S(a) = a
   \langle 2 \rangle 1. Assume: for a contradiction sup S(a) < a
   \langle 2 \rangle 2. Pick b \in B such that \sup S(a) < b < a
   \langle 2 \rangle 3. \ b \in S(a)
   \langle 2 \rangle 4. Q.E.D.
      PROOF: This contradicts the fact that \sup S(a) < b.
\langle 1 \rangle 6. For all b \in S(a) we have b \leq \theta(a)
   \langle 2 \rangle 1. Let: b \in S(a)
   \langle 2 \rangle 2. b < a
   \langle 2 \rangle 3. \ \theta(b) \leq \theta(a)
   \langle 2 \rangle 4. \ b \leq \theta(a)
      Proof: \theta(b) = b
\langle 1 \rangle 7. \ a \leq \theta(a)
  PROOF: Since a = \sup S(a) and \theta(a) is an upper bound for S(a).
\langle 1 \rangle 8. \ a \not< \theta(a)
   \langle 2 \rangle 1. Assume: for a contradiction a < \theta(a).
   \langle 2 \rangle 2. Pick b \in B such that a < b < \theta(a)
   \langle 2 \rangle 3. \theta(a) \leq \theta(b) = b
   \langle 2 \rangle 4. Q.E.D.
      PROOF: This contradicts the fact that b < \theta(a).
\langle 1 \rangle 9. \ \theta(a) = a
```

**Theorem 5.1.21** (Z). Let A and P be complete posets with no least or greatest element. Let B be dense in A and Q be dense in P. Every order isomorphism  $\phi: B \cong Q$  extends uniquely to an order isomorphism  $A \cong P$ .

#### Proof:

```
A ROOF: \langle 1 \rangle1. For a \in A, let S(a) = \{b \in B \mid b < a\}. \langle 1 \rangle2. Define \overline{\phi}: A \to P by \overline{\phi}(a) = \sup \phi(S(a)). \langle 2 \rangle1. \phi(S(a)) is nonempty. \langle 3 \rangle1. PICK a_1 < a
PROOF: Since a is not least. \langle 3 \rangle2. PICK b \in B such that a_1 < b < a. \langle 3 \rangle3. \phi(b) \in \phi(S(a)) \langle 2 \rangle2. \phi(S(a)) is bounded above. \langle 3 \rangle1. PICK a_2 > a
PROOF: Since a is not greatest.
```

 $\langle 3 \rangle 2$ . Pick  $b \in B$  such that  $a < b < a_2$ 

```
\langle 3 \rangle 3. \phi(b) is an upper bound for \phi(S(a)).
\langle 1 \rangle 3. \overline{\phi} is monotone.
   PROOF: If a \leq a' then S(a) \subseteq S(a') and so \overline{\phi}(a) \leq \overline{\phi}(a').
\langle 1 \rangle 4. \phi extends \phi.
   \langle 2 \rangle 1. Let: b \in B
             PROVE: \phi(b) = \sup \phi(S(b))
    \langle 2 \rangle 2. \phi(b) is an upper bound for \phi(S(b))
   \langle 2 \rangle 3. Let: u be any upper bound for \phi(S(b))
             Prove: \phi(b) < u
   \langle 2 \rangle 4. Assume: for a contradiction u < \phi(b)
   \langle 2 \rangle5. Pick q \in Q such that u < q < \phi(b)
   \langle 2 \rangle 6. Pick b' \in B such that \phi(b') = q
    \langle 2 \rangle 7. \ b' < b
    \langle 2 \rangle 8. \ b' \in S(b)
    \langle 2 \rangle 9. \ \ q = \phi(b') \leq u
   \langle 2 \rangle 10. Q.E.D.
       Proof: This is a contradiction.
\langle 1 \rangle 5. Let: \psi = \phi^{-1}
(1)6. Let: \overline{\psi}: P \to A be the function \overline{\psi}(p) = \sup\{\psi(q) \mid q \in Q, q < p\}
\langle 1 \rangle 7. \overline{\psi} is monotone and extends \psi
   PROOF: Similar.
\langle 1 \rangle 8. \overline{\psi} \circ \overline{\phi} : A \to A is monotone and the identity on B.
\langle 1 \rangle 9. \ \overline{\psi} \circ \overline{\phi} = \mathrm{id}_A
   Proof: Proposition 5.1.20.
\langle 1 \rangle 10. \ \overline{\phi} \circ \overline{\psi} = \mathrm{id}_B
   Proof: Proposition 5.1.20.
\langle 1 \rangle 11. If \phi^* : A \cong P is any order isomorphism that extends \phi then \phi^* = \overline{\phi}.
    \langle 2 \rangle 1. Let: a \in A
             PROVE: \phi^*(a) = \sup \phi(S(a))
    \langle 2 \rangle 2. \phi^*(a) is an upper bound for \phi(S(a))
   \langle 2 \rangle 3. Let: u be any upper bound for \phi(S(a))
             PROVE: \phi^*(a) \le u
    \langle 2 \rangle 4. Assume: for a contradiction u < \phi^*(a)
    \langle 2 \rangle5. Pick q \in Q such that u < q < \phi^*(a)
    \langle 2 \rangle 6. Pick b \in B such that q = \phi(b)
    \langle 2 \rangle 7. \ b < a
    \langle 2 \rangle 8. \ b \in S(a)
   \langle 2 \rangle 9. \ q = \phi(b) \leq u
   \langle 2 \rangle 10. Q.E.D.
       PROOF: This is a contradiction.
```

**Definition 5.1.22** (Initial Segment). Let  $\leq$  be a partial order on **A** and  $t \in A$ . The *initial segment* up to t is the class

$$\operatorname{seg} t := \{ x \in \mathbf{A} \mid x < t \} .$$

**Definition 5.1.23** (Lexicographic Ordering). Let **R** be a partial order on **A** and **S** a partial order on **B**. The *lexicographic ordering*  $\leq$  on **A**  $\times$  **B** is defined by:

$$(a,b) \le (a',b') \Leftrightarrow (a\mathbf{R}a' \wedge a \ne a') \vee (a = a' \wedge b\mathbf{S}b')$$
.

**Proposition Schema 5.1.24** (Z). For any classes A, B, R and S, the following is a theorem:

If **R** is a partial order on **A** and **S** is a partial order on **B** then the lexicographic ordering on  $\mathbf{A} \times \mathbf{B}$  is a partial order.

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $\leq$  be the lexicographic ordering on  $\mathbf{A} \times \mathbf{B}$
- $\langle 1 \rangle 2. \leq \text{is reflexive.}$

PROOF: For any  $a \in \mathbf{A}$  and  $b \in \mathbf{B}$  we have a = a and  $b\mathbf{S}b$ , so  $(a, b) \leq (a, b)$ .

- $\langle 1 \rangle 3. \leq \text{is antisymmetric.}$ 
  - (2)1. Assume:  $(a,b) \le (a',b')$  and  $(a',b') \le (a,b)$
  - $\langle 2 \rangle 2$ .  $(a\mathbf{R}a' \wedge a \neq a') \vee (a = a' \wedge b\mathbf{S}b')$
  - $\langle 2 \rangle 3$ .  $(a' \mathbf{R} a \wedge a' \neq a) \vee (a' = a \wedge b \mathbf{S} b')$
  - $\langle 2 \rangle 4$ . Case: a = a'

PROOF: Then  $b\mathbf{S}b'$  and  $b'\mathbf{S}b$  hence b=b' and (a,b)=(a',b').

 $\langle 2 \rangle$ 5. Case:  $a \neq a'$ 

PROOF: Then  $a\mathbf{R}a'$  and  $a'\mathbf{R}a$  hence a=a' which is a contradiction.

- $\langle 1 \rangle 4$ .  $\leq$  is transitive.
  - $\langle 2 \rangle 1$ . Assume:  $(a_1, b_1) \leq (a_2, b_2) \leq (a_3, b_3)$
  - $\langle 2 \rangle 2$ .  $(a_1 \mathbf{R} a_2 \wedge a_1 \neq a_2) \vee (a_1 = a_2 \wedge b_1 \mathbf{S} b_2)$
  - $\langle 2 \rangle 3. \ (a_2 \mathbf{R} a_3 \wedge a_2 \neq a_3) \vee (a_2 = a_3 \wedge b_2 \mathbf{S} b_3)$
  - $\langle 2 \rangle 4$ . Case:  $a_1 \mathbf{R} a_2, a_1 \neq a_2, a_2 \mathbf{R} a_3, a_2 \neq a_3$ 
    - $\langle 3 \rangle 1. \ a_1 \mathbf{R} a_3$

PROOF: Since  $\mathbf{R}$  is transitive.

 $\langle 3 \rangle 2$ .  $a_1 \neq a_3$ 

PROOF: If  $a_1 = a_3$  then  $a_1 \mathbf{R} a_2$  and  $a_2 \mathbf{R} a_1$  so  $a_1 = a_2$  which is a contradiction.

 $\langle 2 \rangle 5$ . Case:  $a_1 \mathbf{R} a_2, a_1 \neq a_2, a_2 = a_3, b_2 \mathbf{S} b_3$ 

PROOF: Then  $a_1 \mathbf{R} a_3$  and  $a_1 \neq a_3$ .

 $\langle 2 \rangle 6$ . Case:  $a_1 = a_2, b_1 \mathbf{S} b_2, a_2 \mathbf{R} a_3, a_2 \neq a_3$ 

PROOF: Then  $a_1 \mathbf{R} a_3$  and  $a_1 \neq a_3$ .

 $\langle 2 \rangle 7$ . Case:  $a_1 = a_2, b_1 \mathbf{S} b_2, a_2 = a_3, b_2 \mathbf{S} b_3$ 

PROOF: Then  $a_1 = a_3$  and  $b_1 \mathbf{S} b_3$ .

## 5.2 Linear Orders

**Definition 5.2.1** (Linear Ordering). Let **A** be a class. A *linear ordering* or *total ordering* on **A** is a partial ordering  $\leq$  on **A** that is *total*, i.e.

$$\forall x,y \in \mathbf{A}.x \leq y \vee y \leq x$$

We often use the symbol < for a linear ordering, and then write x < y for  $(x,y) \in <$ .

**Proposition Schema 5.2.2** (Trichotomy (Z)). For any classes **A** and  $\leq$ , the following is a theorem:

Assume  $\leq$  be a linear ordering on **A**. For any  $x, y \in \mathbf{A}$ , exactly one of x < y, x = y, y < x holds.

Proof: Immediate from definitions.  $\Box$ 

**Proposition Schema 5.2.3** (Z). For any classes A and <, the following is a theorem:

Let < be a transitive relation on  $\mathbf{A}$  that satisfies trichotomy. Define  $\leq$  on  $\mathbf{A}$  by  $x \leq y$  iff x < y or x = y. Then  $\leq$  is a linear ordering on  $\mathbf{A}$  and x < y iff  $x \leq y$  and  $x \neq y$ .

#### Proof:

 $\langle 1 \rangle 1$ . < is reflexive.

PROOF: By definition we have  $\forall x \in \mathbf{A}.x \leq x$ .

- $\langle 1 \rangle 2$ .  $\leq$  is antisymmetric.
  - $\langle 2 \rangle 1$ . Assume:  $x \leq y$  and  $y \leq x$
  - $\langle 2 \rangle 2$ . x < y or x = y
  - $\langle 2 \rangle 3$ . y < x or y = x
  - $\langle 2 \rangle 4$ . We cannot have x < y and y < x

PROOF: Trichotomy.

- $\langle 2 \rangle 5. \ x = y$
- $\langle 1 \rangle 3. \leq \text{is transitive.}$ 
  - $\langle 2 \rangle 1$ . Assume:  $x \leq y$  and  $y \leq z$
  - $\langle 2 \rangle 2$ . x < y or x = y
  - $\langle 2 \rangle 3$ . y < z or y = z
  - $\langle 2 \rangle 4$ . Case: x < y and y < z

PROOF: Then x < z by transitivity, so  $x \le z$ .

 $\langle 2 \rangle 5$ . Case: x = y

PROOF: Then we have  $y \leq z$  and so  $x \leq z$ .

 $\langle 2 \rangle 6$ . Case: y = z

PROOF: Then we have  $x \leq y$  and so  $x \leq z$ .

 $\langle 1 \rangle 4. \leq \text{is total.}$ 

PROOF: Immediate from trichotomy.

**Proposition Schema 5.2.4** (Z). For any classes **A** and **R**, the following is a theorem:

If  $\mathbf{R}$  is a linear ordering on  $\mathbf{A}$  then  $\mathbf{R}^{-1}$  is also a linear ordering on  $\mathbf{A}$ .

#### PROOF

 $\langle 1 \rangle 1$ .  $\mathbf{R}^{-1}$  is a partial order on  $\mathbf{A}$ .

Proof: Proposition 5.1.2.

 $\langle 1 \rangle 2$ .  $\mathbf{R}^{-1}$  is total.

```
\langle 2 \rangle 1. Let: x, y \in \mathbf{A}

\langle 2 \rangle 2. x\mathbf{R}y or y\mathbf{R}x.

\langle 2 \rangle 3. y\mathbf{R}^{-1}x or x\mathbf{R}^{-1}y.
```

**Proposition Schema 5.2.5** (Z). For any classes **A**, **B**, **F**, **R**, **S**, the following is a theorem:

Assume **R** is a linear order on **A**, **S** is a partial order on **B**, and **F** :  $\mathbf{A} \to \mathbf{B}$ . If **F** is strictly monotone then it is injective.

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $x, y \in \mathbf{A}$
- $\langle 1 \rangle 2$ . Assume:  $x \neq y$

PROVE:  $\mathbf{F}(x) \neq \mathbf{F}(y)$ 

 $\langle 1 \rangle 3$ . Assume: w.l.o.g.  $x \mathbf{R} y$ 

PROOF:  $\mathbf{R}$  is total.

 $\langle 1 \rangle 4$ .  $\mathbf{F}(x)\mathbf{SF}(y)$  and  $\mathbf{F}(x) \neq \mathbf{F}(y)$ 

PROOF: **F** is strictly monotone.

**Proposition Schema 5.2.6** (Z). For any classes A, B,  $\leq$ ,  $\preccurlyeq$  and F, the following is a theorem:

Assume  $\leq$  is a linear order on  $\mathbf{A}$  and  $\preccurlyeq$  is a linear order on  $\mathbf{B}$ . Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{F}$  is strictly monotone. For all  $x, y \in \mathbf{A}$ , if  $\mathbf{F}(x) \prec \mathbf{F}(y)$  then x < y.

#### Proof:

 $\langle 1 \rangle 1$ .  $\mathbf{F}(x) \neq \mathbf{F}(y)$  and  $\mathbf{F}(y) \not\prec \mathbf{F}(x)$ 

PROOF: Trichotomy.

 $\langle 1 \rangle 2$ .  $x \neq y$  and  $y \not< x$ 

Proof:  $\mathbf{F}$  is strictly monotone.

 $\langle 1 \rangle 3. \ x < y$ 

Proof: Trichotomy.

**Corollary Schema 5.2.6.1** (Z). For any classes A, B,  $\leq$ ,  $\preccurlyeq$  and F, the following is a theorem:

Assume  $\leq$  is a linear order on  $\mathbf{A}$  and  $\preccurlyeq$  is a linear order on  $\mathbf{B}$ . Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{F}$  is strictly monotone. Then  $\mathbf{F}$  is an order isomorphism.

**Proposition Schema 5.2.7** (Z). For any classes **A**, **B**, **F** and **S**, the following is a theorem:

Assume **S** is a linear order on **B** and **F**:  $\mathbf{A} \rightarrow \mathbf{B}$ . Define **R** on **A** by  $x\mathbf{R}y$  if and only if  $\mathbf{F}(x)\mathbf{SF}(y)$ . Then **R** is a linear order on **A**.

### Proof:

 $\langle 1 \rangle 1$ . **R** is a partial order on **A**.

Proof: Proposition 5.1.3.

```
\langle 1 \rangle 2. R is total. PROOF: For all x, y \in \mathbf{A} we have \mathbf{F}(x)\mathbf{SF}(y) or \mathbf{F}(y)\mathbf{SF}(x).
```

Corollary Schema 5.2.7.1 (Z). For any classes A, B and R, the following is a theorem:

Assume **R** be a linear order on **A** and **B**  $\subseteq$  **A**. Then **R**  $\cap$  **B**<sup>2</sup> is a linear order on **B**.

**Proposition Schema 5.2.8** (Z). For any classes **A**, **B**, **R** and **S**, the following is a theorem:

Assume  $\mathbf{R}$  is a linear order on  $\mathbf{A}$  and  $\mathbf{S}$  is a linear order on  $\mathbf{B}$ . Then the lexicographic ordering is a linear order on  $\mathbf{A} \times \mathbf{B}$ .

```
Proof:
```

```
\begin{array}{l} \langle 1 \rangle 1. \text{ Let: } \leq \text{ be the lexicographic order on } \mathbf{A} \times \mathbf{B} \\ \langle 1 \rangle 2. \leq \text{ is a partial order.} \\ \text{Proof: Proposition 5.1.24.} \\ \langle 1 \rangle 3. \leq \text{ is total.} \\ \langle 2 \rangle 1. \text{ Let: } a, a' \in \mathbf{A} \text{ and } b, b' \in \mathbf{B} \\ \langle 2 \rangle 2. \text{ Case: } a\mathbf{R}a' \text{ and } a \neq a' \\ \text{Proof: Then } (a,b) \leq (a',b'). \\ \langle 2 \rangle 3. \text{ Case: } a = a' \\ \text{Proof: We have } b\mathbf{S}b' \text{ or } b'\mathbf{S}b, \text{ so } (a,b) \leq (a',b') \text{ or } (a',b') \leq (a,b). \\ \langle 2 \rangle 4. \text{ Case: } a'\mathbf{R}a \text{ and } a \neq a' \\ \text{Proof: Then } (a',b') \leq (a,b). \\ \end{array}
```

# 5.3 Well Orderings

**Definition 5.3.1** (Well Ordering). A well ordering on a class A is a well-founded linear ordering on A.

**Proposition Schema 5.3.2.** For any classes A, B, F and S, the following is a theorem:

Assume **S** well orders **B** and **F** :  $\mathbf{A} \rightarrow \mathbf{B}$ . Define **R** on **A** by  $x\mathbf{R}y$  if and only if  $\mathbf{F}(x)\mathbf{SF}(y)$ . Then **R** well orders **A**.

#### Proof:

```
\langle 1 \rangle 1. R linearly orders A.
```

Proof: Proposition 5.2.7.

- $\langle 1 \rangle 2$ . For all  $t \in \mathbf{A}$  we have  $\{x \in \mathbf{A} \mid x\mathbf{R}t \land x \neq t\}$  is a set.
  - $\langle 2 \rangle 1$ . Let:  $t \in \mathbf{A}$
  - $\langle 2 \rangle 2$ . Let:  $S = \{ y \in \mathbf{B} \mid y\mathbf{SF}(t) \land y \neq \mathbf{F}(t) \}$
  - $\langle 2 \rangle 3$ . Let: P(x,y) be the property  $\mathbf{F}(y) = x$
  - $\langle 2 \rangle 4$ . For all  $x \in S$  there exists at most one y such that P(x,y)

Proof:  $\mathbf{F}$  is injective.

```
\langle 2 \rangle5. Let: T = \{y \mid \exists x \in S.P(x,y)\}
Proof: Axiom of Replacement.
\langle 2 \rangle6. T = \{x \in \mathbf{A} \mid x\mathbf{R}t \land x \neq t\}
\langle 1 \rangle3. Every nonempty subset of \mathbf{A} has a least element.
\langle 2 \rangle1. Let: S be a nonempty subset of \mathbf{A}.
\langle 2 \rangle2. Let: y be the least element of \mathbf{F}(S).
\langle 2 \rangle3. Pick x \in S such that \mathbf{F}(x) = y.
\langle 2 \rangle4. x is least in S.
```

Corollary Schema 5.3.2.1. For any classes A, B and R, the following is a theorem:

If **R** well orders **B** and  $\mathbf{A} \subseteq \mathbf{B}$  then  $\mathbf{R} \cap \mathbf{A}^2$  well orders **A**.

**Proposition 5.3.3.** For any well ordered sets A and B, the lexicographic order well orders  $A \times B$ .

#### Proof:

 $\langle 1 \rangle 1$ .  $A \times B$  is linearly ordered.

Proof: Proposition 5.2.8.

- $\langle 1 \rangle 2$ . Every nonempty subset of  $A \times B$  has a least element.
  - $\langle 2 \rangle 1$ . Let: S be a nonempty subset of  $A \times B$ .
  - $\langle 2 \rangle 2$ . Let: a be the least element of  $\{x \in A \mid \exists y \in B.(x,y) \in S\}$ .
  - $\langle 2 \rangle 3$ . Let: b be the least element of  $\{ y \in B \mid (a, y) \in S \}$ .
- $\langle 2 \rangle 4$ . (a,b) is least in S.

**Definition 5.3.4** (End Extension). Let A and B be well ordered sets. Then B is an *end extension* of A iff  $A \subseteq B$  and:

- Whenever  $x, y \in A$  then  $x \leq_A y$  iff  $x \leq_B y$ .
- Whenever  $x \in A$  and  $y \in B A$  then x < y.

**Theorem 5.3.5.** Let  $\leq$  be a linear ordering on A. Assume that, for any  $B \subseteq A$  such that  $\forall t \in A$ . seg  $t \subseteq B \Rightarrow t \in B$ , we have B = A. Then  $\leq$  is a well ordering on A.

## Proof:

- $\langle 1 \rangle 1$ . Let:  $C \subseteq A$  be nonempty.
- $\langle 1 \rangle 2$ . Let:  $B = \{ t \in A \mid \forall x \in C . t < x \}$
- $\langle 1 \rangle 3. \ B \cap C = \emptyset$
- $\langle 1 \rangle 4. \ B \neq A$
- $\langle 1 \rangle 5$ . Pick  $t \in A$  such that  $\operatorname{seg} t \subseteq B$  and  $t \notin B$
- $\langle 1 \rangle 6$ . t is least in C.

**Proposition Schema 5.3.6.** For any classes A, B, F, G,  $\leq$  and  $\leq$ , the following is a theorem:

Assume  $\leq$  well orders  $\mathbf{A}$  and  $\leq$  well orders  $\mathbf{B}$ . Assume  $\mathbf{F}$  and  $\mathbf{G}$  are order isomorphisms between  $\mathbf{A}$  and  $\mathbf{B}$ . Then  $\mathbf{F} = \mathbf{G}$ .

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Proof:

$$\langle 1 \rangle 1$$
. For all  $x \in \mathbf{A}$ , if  $\forall t < x.\mathbf{F}(t) = \mathbf{G}(t)$ , then  $\mathbf{F}(x) = \mathbf{G}(x)$ 

 $\langle 2 \rangle 1$ . Let:  $x \in \mathbf{A}$ 

 $\langle 2 \rangle 2$ . Assume:  $\forall t < x. \mathbf{F}(t) = \mathbf{G}(t)$ 

 $\langle 2 \rangle 3$ .  $\mathbf{F}(\operatorname{seg} x) = \mathbf{G}(\operatorname{seg} x)$ 

 $\langle 2 \rangle 4$ .  $\mathbf{F}(x)$  is the least element of  $\mathbf{B} - \mathbf{F}(\operatorname{seg} x)$ 

 $\langle 2 \rangle$ 5.  $\mathbf{G}(x)$  is the least element of  $\mathbf{B} - \mathbf{G}(\operatorname{seg} x)$ 

 $\langle 2 \rangle 6. \ \mathbf{F}(x) = \mathbf{G}(x)$ 

 $\langle 1 \rangle 2. \ \forall x \in \mathbf{A}.\mathbf{F}(x) = \mathbf{G}(x)$ 

PROOF: Transfinite induction.

**Theorem 5.3.7.** Let A and B be well ordered sets. Then one of the following holds:  $A \cong B$ ; there exists  $b \in B$  such that  $A \cong \text{seg } b$ ; there exists  $a \in A$  such that  $\text{seg } a \cong B$ .

Proof:

 $\langle 1 \rangle 1$ . PICK e that is not in A or B.

 $\langle 1 \rangle 2$ . Let:  $F: A \to B \cup \{e\}$  be the function defined by transfinite recursion thus:

$$F(t) = \begin{cases} \text{the least element of } B - F(\sec t) & \text{if } B - F(\sec t) \neq \emptyset \\ e & \text{if } B - F(\sec t) = \emptyset \end{cases}$$

 $\langle 1 \rangle 3$ . Case:  $e \in \operatorname{ran} F$ 

 $\langle 2 \rangle 1$ . Let: t be least such that F(t) = e

 $\langle 2 \rangle 2$ .  $F \upharpoonright \operatorname{seg} t : \operatorname{seg} t \cong B$ 

 $\langle 1 \rangle 4$ . Case: ran F = B

PROOF: We have  $F:A\cong B$ 

 $\langle 1 \rangle$ 5. Case: ran  $F \subseteq B$ 

 $\langle 2 \rangle 1$ . Let: b be the least element of  $B - \operatorname{ran} F$ 

 $\langle 2 \rangle 2$ .  $F: A \cong \text{seg } b$ 

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# Chapter 6

# **Ordinal Numbers**

## 6.1 Ordinals

**Definition 6.1.1** (Ordinal Number). An *ordinal (number)* is a transitive set  $\alpha$  that is *well-ordered by*  $\in$ ; that is, such that  $\{(x,y) \in \alpha^2 \mid x \in y \lor x = y\}$  well orders  $\alpha$ .

Given  $x, y \in \alpha$ , we write x < y iff  $x \in y$ , and  $x \le y$  iff  $x \in y$  or x = y.

Let **On** be the class of ordinal numbers. For  $\alpha, \beta \in$  **On**, we write  $\alpha < \beta$  iff  $\alpha \in \beta$ , and  $\alpha \leq \beta$  iff  $\alpha < \beta$  or  $\alpha = \beta$ .

**Proposition 6.1.2.** For any ordinal numbers  $\alpha$  and  $\beta$ , if  $\alpha \cong \beta$  then  $\alpha = \beta$ .

```
Proof:
\langle 1 \rangle 1. Let: f : \alpha \cong \beta
\langle 1 \rangle 2. For all x \in \alpha, if \forall t < x. f(t) = t then f(x) = x
    \langle 2 \rangle 1. \ f(x) \subseteq x
        \langle 3 \rangle 1. Let: y \in f(x)
        \langle 3 \rangle 2. \ y \in \beta
        \langle 3 \rangle 3. Pick t \in \alpha such that f(t) = y
            PROOF: f is surjective.
        \langle 3 \rangle 4. \ f(t) \in f(x)
        \langle 3 \rangle 5. \ t \in x
            PROOF: Since f is an order isomorphism.
        \langle 3 \rangle 6. f(t) = t
            Proof: Induction hypothesis.
        \langle 3 \rangle 7. \ y = t
        \langle 3 \rangle 8. \ y \in x
    \langle 2 \rangle 2. x \subseteq f(x)
        \langle 3 \rangle 1. Let: t \in x
        \langle 3 \rangle 2. \ f(t) \in f(x)
        \langle 3 \rangle 3. \ f(t) = t
        \langle 3 \rangle 4. \ t \in f(x)
```

```
\langle 1 \rangle 3. \ \forall x \in \alpha. f(x) = x PROOF: Transfinite induction. \langle 1 \rangle 4. \ \alpha = \beta PROOF: Since \beta = \{f(t) \mid t \in \alpha\} = \{t \mid t \in \alpha\} = \alpha. \Box

Theorem 6.1.3. Every well-ordered set is isomorphic to a unique ordinal. PROOF: \langle 1 \rangle 1. For any well-ordered set A, there exists an ordinal \alpha such that A \cong \alpha. \langle 2 \rangle 1. Let: A be a well-ordered set. \langle 2 \rangle 2. Define the function E on A by transfinite recursion thus: E(t) = \{E(x) \mid x < t\} (t \in A).
```

 $\langle 4 \rangle$ 2. PICK  $t \in A$  such that y = E(t) $\langle 4 \rangle$ 3.  $x \in E(t) = \{E(s) \mid s < t\}$  $\langle 4 \rangle$ 4. PICK s < t such that x = E(s)

 $\langle 4 \rangle 1$ . Let:  $\langle = \{(x,y) \in \alpha \mid x \in y\}$ 

 $\langle 5 \rangle 1$ . Let:  $x, y, z \in \alpha$  with  $x \in y \in z$  $\langle 5 \rangle 2$ . Pick  $t \in A$  such that z = E(t)

 $\langle 5 \rangle 3$ . Pick  $s \in A$  such that s < t and y = E(s)  $\langle 5 \rangle 4$ . Pick  $r \in A$  such that r < s and x = E(r)

 $\langle 5 \rangle 2$ . PICK  $s, t \in A$  such that E(s) = x and E(t) = y

 $\langle 5 \rangle 3$ . Exactly one of s < t, s = t, t < s holds.

PROOF: Axiom of Regularity.

PROOF: Axiom of Regularity.

 $\langle 4 \rangle$ 5. Every nonempty subset of  $\alpha$  has a least element.

 $\langle 2 \rangle 3$ . Let:  $\alpha = \{ E(x) \mid x \in A \}$ 

 $\langle 3 \rangle 1$ .  $\alpha$  is a transitive set.  $\langle 4 \rangle 1$ . Let:  $x \in y \in \alpha$ 

 $\langle 3 \rangle 2$ .  $\alpha$  is well-ordered by  $\in$ .

 $\langle 4 \rangle 3$ . < satisfies trichotomy.  $\langle 5 \rangle 1$ . Let:  $x, y \in \alpha$ 

 $\langle 6 \rangle 2$ .  $x \neq y$  and  $y \notin x$ 

 $\langle 6 \rangle 2$ .  $x \notin y$  and  $y \notin x$ 

PROOF: Similar to  $\langle 5 \rangle 4$ .  $\langle 4 \rangle 4$ .  $\leq$  is a linear order on  $\alpha$ . PROOF: Proposition 5.2.3.

 $\langle 5 \rangle 4$ . Case: s < t $\langle 6 \rangle 1$ .  $x \in y$ 

 $\langle 5 \rangle 5$ . Case: s = t $\langle 6 \rangle 1$ . x = y

 $\langle 5 \rangle 6$ . Case: t < s

 $\langle 4 \rangle 2$ . < is transitive.

 $\langle 5 \rangle 5. \ r < t$  $\langle 5 \rangle 6. \ x \in z$ 

 $\langle 2 \rangle 4$ .  $\alpha$  is an ordinal.

 $\langle 4 \rangle 5. \ x \in \alpha$ 

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```
\langle 5 \rangle 1. Let: S be a nonempty subset of \alpha
              \langle 5 \rangle 2. Let: T = \{ x \in A \mid E(x) \in S \}
              \langle 5 \rangle 3. Let: t be the least element of T.
                       PROVE: E(t) is least in S
              \langle 5 \rangle 4. Let: y \in S
              \langle 5 \rangle 5. Pick s \in T such that E(s) = y
              \langle 5 \rangle 6. \ t \leq s
              \langle 5 \rangle 7. \ x \leq y
   \langle 2 \rangle5. E is surjective.
       PROOF: By definition of \alpha.
   \langle 2 \rangle 6. E is strictly monotone.
       PROOF: If s < t then E(s) \in E(t) by definition of E(t).
   \langle 2 \rangle7. Q.E.D.
       Proof: Corollary 5.2.6.1.
\langle 1 \rangle 2. For any ordinals \alpha and \beta, if \alpha \cong \beta then \alpha = \beta.
   Proof: Proposition 6.1.2.
Proposition 6.1.4. The class On is a transitive class. That is, every element
of an ordinal is an ordinal.
Proof:
\langle 1 \rangle 1. Let: \alpha be an ordinal.
\langle 1 \rangle 2. Let: \beta \in \alpha
\langle 1 \rangle 3. \beta is a transitive set.
   \langle 2 \rangle 1. Let: x \in y \in \beta
   \langle 2 \rangle 2. \ y \in \alpha
       PROOF: \alpha is transitive.
   \langle 2 \rangle 3. \ x \in \alpha
       PROOF: \alpha is transitive.
   \langle 2 \rangle 4. \ x \in \beta
       PROOF: Since \{(x,y) \in \alpha^2 \mid x \in y\} is transitive.
\langle 1 \rangle 4. \beta is well ordered by \in.
   Proof: By Corollary 5.3.2.1.
Proposition 6.1.5. Given two ordinal numbers \alpha, \beta, exactly one of \alpha \in \beta,
\alpha = \beta, \ \beta \in \alpha \ holds.
Proof:
\langle 1 \rangle 1. At most one holds.
   PROOF: Since every ordinal is a transitive set and we never have \alpha \in \alpha.
\langle 1 \rangle 2. At least one holds.
   \langle 2 \rangle 1. Either \alpha \cong \beta or \exists t \in \beta . \alpha \cong \text{seg } t or \exists t \in \alpha . \text{seg } t \cong \beta .
   \langle 2 \rangle 2. Case: \alpha \cong \beta
       PROOF: Then \alpha = \beta by Proposition 6.1.2.
   \langle 2 \rangle 3. Case: There exists t \in \beta such that \alpha \cong \operatorname{seg} t
```

```
\langle 3 \rangle 1. t is an ordinal number.
          Proof: Proposition 6.1.4.
       \langle 3 \rangle 2. t = \sec t
          \langle 4 \rangle 1. t \subseteq \operatorname{seg} t
              \langle 5 \rangle 1. Let: s \in t
              \langle 5 \rangle 2. \ s \in \beta
                 PROOF: \beta is a transitive set.
              \langle 5 \rangle 3. \ s \in \operatorname{seg} t
          \langle 4 \rangle 2. seg t \subseteq t
              PROOF: Immediate from definitions.
       \langle 3 \rangle 3. \ \alpha = t
          Proof: Proposition 6.1.2.
       \langle 3 \rangle 4. \ \alpha \in \beta
    \langle 2 \rangle 4. Case: There exists t \in \alpha such that \operatorname{seg} t \cong \beta
       PROOF: \beta \in \alpha similarly.
Proposition 6.1.6. Any nonempty set S of ordinal numbers has a least ele-
ment.
Proof:
\langle 1 \rangle 1. Ріск \beta \in S
\langle 1 \rangle 2. Case: \beta \cap S = \emptyset
   PROOF: Then \beta is least in S.
\langle 1 \rangle 3. Case: \beta \cap S \neq \emptyset
   PROOF: The least element of \beta \cap S is least in S.
Theorem 6.1.7. The class On is well ordered by \in.
Proof:
\langle 1 \rangle 1. Let: \mathbf{E} = \{ (x, y) \in \mathbf{On}^2 \mid x \in y \}
\langle 1 \rangle 2. E is transitive.
   PROOF: If \alpha \in \beta \in \gamma then \alpha \in \gamma because every ordinal is a transitive set.
\langle 1 \rangle 3. E satisfies trichotomy.
   Proof: Proposition 6.1.5.
\langle 1 \rangle 4. E linearly orders On.
   Proof: Proposition 5.2.3.
\langle 1 \rangle 5. For any \alpha \in \mathbf{On}, the class \{ \beta \in \mathbf{On} \mid \beta \mathbf{E} \alpha \} is a set.
   PROOF: It is equal to \alpha.
\langle 1 \rangle 6. Every nonempty subset of On has an E-least element.
   Proof: Proposition 6.1.6.
```

Corollary 6.1.7.1 (Burali-Forti Paradox). The class On is a proper class.

PROOF: If it were a set, it would be a transitive set well-ordered by  $\in$ , and hence a member of itself, contradicting the Axiom of Regularity.

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**Proposition 6.1.8.** Any transitive set of ordinal numbers is an ordinal number.

PROOF: It is well-ordered by  $\in$  by Corollary 5.3.2.1 and Theorem 6.1.7.  $\square$ 

**Proposition 6.1.9.**  $\emptyset$  is an ordinal number.

PROOF: Vacuously, it is a transitive set well-ordered by  $\in$ .

**Definition 6.1.10.** We define  $0 = \emptyset$ .

**Proposition 6.1.11.** If A is a set of ordinal numbers then  $\bigcup A$  is an ordinal number.

### Proof:

 $\langle 1 \rangle 1$ .  $\bigcup A$  is a transitive set.

PROOF: Proposition 1.6.3.

 $\langle 1 \rangle 2$ .  $\bigcup A$  is a set of ordinals.

Proof: Proposition 6.1.4.

Corollary 6.1.11.1. The poset On is complete.

PROOF: For any nonempty set A of ordinals,  $\bigcup A$  is its supremum.  $\square$ 

**Proposition 6.1.12.** Let  $\alpha$  be an ordinal and  $S \subseteq \alpha$ . Then S is well-ordered by  $\in$  and the ordinal of  $(S, \in)$  is  $\leq \alpha$ .

#### Proof:

- $\langle 1 \rangle 1$ . S is well ordered by  $\in$ .
- $\langle 1 \rangle 2$ . Let:  $\beta$  be the ordinal of  $(S, \in)$
- $\langle 1 \rangle 3$ . Let:  $E: S \approx \beta$  be the unique isomorphism.
- $\langle 1 \rangle 4. \ \forall \gamma \in S.E(\gamma) \leq \gamma$ 
  - $\langle 2 \rangle 1$ . Let:  $\gamma \in S$
  - $\langle 2 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall \delta < \gamma . E(\delta) \leq \delta$
  - $\langle 2 \rangle 3$ .  $E(\gamma)$  is the least element of  $\beta$  that is greater than  $E(\delta)$  for all  $\delta < \gamma$
  - $\langle 2 \rangle 4$ .  $\gamma$  is greater than  $E(\delta)$  for all  $\delta < \gamma$
  - $\langle 2 \rangle 5$ .  $E(\gamma) \leq \gamma$
- $\langle 1 \rangle 5$ .  $\beta < \alpha$ 
  - $\langle 2 \rangle 1. \ \forall \gamma < \beta. \gamma < \alpha$ 
    - $\langle 3 \rangle 1$ . Let:  $\gamma < \beta$
    - $\langle 3 \rangle 2$ . Pick  $\delta \in S$  such that  $E(\delta) = \gamma$
    - $\langle 3 \rangle 3. \ \gamma = E(\delta) \le \delta < \alpha$

**Proposition 6.1.13.** Let  $\alpha$  be a set. Then the following are equivalent.

- 1.  $\alpha$  is an ordinal.
- 2.  $\alpha$  is a transitive set and, for all  $x, y \in \alpha$ , either x = y or  $x \in y$  or  $y \in x$ .
- 3.  $\alpha$  is a transitive set of transitive sets.

```
Proof:
```

 $\langle 1 \rangle 1. \ 1 \Rightarrow 2$ 

PROOF: Immediate from definitions.

- $\langle 1 \rangle 2. \ 2 \Rightarrow 3$ 
  - $\langle 2 \rangle 1$ . Assume:  $\alpha$  is a transitive set and, for all  $x, y \in \alpha$ , either x = y or  $x \in y$ or  $y \in x$
  - $\langle 2 \rangle 2$ . Let:  $z \in \alpha$

Prove: z is transitive.

- $\langle 2 \rangle 3$ . Let:  $x \in y \in z$
- $\langle 2 \rangle 4. \ y \in \alpha$
- $\langle 2 \rangle 5. \ x \in \alpha$
- $\langle 2 \rangle 6$ . Either x = z or  $x \in z$  or  $z \in x$
- $\langle 2 \rangle 7. \ x \neq z$

PROOF: We cannot have  $x \in y \in x$  by the Axiom of Regularity.

PROOF: We cannot have  $x \in y \in z \in x$  by the Axiom of Regularity.

- $\langle 1 \rangle 3. \ 3 \Rightarrow 1$ 
  - $\langle 2 \rangle 1$ . Let: x be a transitive set of transitive sets.
  - $\langle 2 \rangle 2$ . Assume: as  $\in$ -induction hypothesis that, for all  $y \in x$ , if y is a transitive set of transitive sets then y is a transitive set of ordinals.
  - $\langle 2 \rangle 3$ . Every element of x is an ordinal.
    - $\langle 3 \rangle 1$ . Let:  $y \in x$
    - $\langle 3 \rangle 2$ . y is transitive.
    - $\langle 3 \rangle 3$ . Every element of y is transitive.

PROOF: Since every element of y is an element of x, because x is transitive.

 $\langle 3 \rangle 4$ . y is an ordinal.

Proof:  $\langle 2 \rangle 2$ 

 $\langle 2 \rangle 4$ . Q.E.D.

Proof: Proposition 6.1.8.

**Lemma 6.1.14.** Let A and B be well-ordered sets. If B is an end extension of A then the ordinal of A is  $\leq$  the ordinal of B.

- $\langle 1 \rangle 1$ . Let:  $\alpha$  be the ordinal of A and  $\beta$  the ordinal of B.
- $\langle 1 \rangle 2$ . Let:  $E_A : A \cong \alpha$  and  $E_B : B \cong \beta$  be the canonical isomorphisms.
- $\langle 1 \rangle 3. \ \forall a \in A.E_A(a) = E_B(a)$ 
  - $\langle 2 \rangle 1$ . Let:  $a \in A$
  - $\langle 2 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall x < a.E_A(x) = E_B(x)$
  - $\langle 2 \rangle 3$ .  $E_A(a)$  is the least ordinal that is greater than  $E_A(x)$  for all x < a
  - $\langle 2 \rangle 4$ .  $E_B(a)$  is the least ordinal that is greater than  $E_B(x)$  for all x < b
  - $\langle 2 \rangle 5. \{ x \in A \mid x <_A a \} = \{ x \in B \mid x <_B a \}$
- $\langle 2 \rangle 6$ .  $E_A(a) = E_B(a)$
- $\langle 1 \rangle 4. \ \alpha \subseteq \beta$
- $\langle 1 \rangle 5. \ \alpha \leq \beta$

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**Lemma 6.1.15.** Let C be a set of well ordered sets such that, for any  $A, B \in C$ , we have that one of A and B is an end extension of the other. Let  $W = \bigcup \mathcal{C}$ under  $x \leq y$  iff there exists  $A \in W$  such that  $x, y \in A$  and  $x \leq y$ . Then W is a well ordered set whose ordinal is the supremum of the ordinals of the members of C.

#### Proof:

- $\langle 1 \rangle 1. \leq \text{is reflexive on } W.$ 
  - $\langle 2 \rangle 1$ . Let:  $x \in W$
  - $\langle 2 \rangle 2$ . PICK  $A \in W$  such that  $x \in A$ .
  - $\langle 2 \rangle 3. \ x \leq x$
- $\langle 1 \rangle 2. \leq \text{is antisymmetric on } W.$ 
  - $\langle 2 \rangle 1$ . Let:  $x, y \in W$
  - $\langle 2 \rangle 2$ . Assume:  $x \leq y$  and  $y \leq x$
  - $\langle 2 \rangle 3$ . PICK  $A \in W$  such that  $x, y \in A$  and  $x \leq_A y$ , and  $B \in W$  such that  $x, y \in B$  and  $y \leq_B x$
  - $\langle 2 \rangle 4$ . Assume: w.l.o.g. B is an end extension of A
  - $\langle 2 \rangle$ 5.  $x \leq_B y$  and  $y \leq_B x$
  - $\langle 2 \rangle 6. \ x = y$
- $\langle 1 \rangle 3. \leq \text{is transitive on } W.$ 
  - $\langle 2 \rangle 1$ . Assume:  $x \leq y \leq z$
  - $\langle 2 \rangle 2$ . PICK  $A, B \in W$  such that  $x \leq_A y$  and  $y \leq_B z$
  - $\langle 2 \rangle 3$ . Case: A is an end extension of B.
    - $\langle 3 \rangle 1$ .  $x \leq_A y$  and  $y \leq_A z$
    - $\langle 3 \rangle 2. \ x \leq_A z$
    - $\langle 3 \rangle 3. \ x \leq z$
  - $\langle 2 \rangle 4$ . Case: B is an end extension of A.

PROOF: Similar.

- $\langle 1 \rangle 4. \leq \text{is total on } W.$ 
  - $\langle 2 \rangle 1$ . Let:  $x, y \in W$
  - $\langle 2 \rangle 2$ . Pick  $A, B \in \mathcal{C}$  such that  $x \in A$  and  $y \in B$
  - $\langle 2 \rangle$ 3. Assume: w.l.o.g. B is an end extension of A
  - $\langle 2 \rangle 4$ .  $x \leq_B y$  or  $y \leq_B x$
  - $\langle 2 \rangle 5$ .  $x \leq_W y$  or  $y \leq_W x$
- $\langle 1 \rangle$ 5. Every nonempty subset of W has a least element.
  - $\langle 2 \rangle 1$ . Let: S be a nonempty subset of W
  - $\langle 2 \rangle 2$ . Pick  $s \in S$
  - $\langle 2 \rangle 3$ . Pick  $A \in \mathcal{C}$  such that  $s \in A$
  - $\langle 2 \rangle 4$ . Let: a be the  $\leq_A$ -least element of  $S \cap A$ Prove: a is least in S
  - $\langle 2 \rangle 5$ . Let:  $x \in S$

Prove:  $a \le x$ 

- $\langle 2 \rangle 6$ . Pick  $B \in \mathcal{C}$  such that  $x \in B$
- $\langle 2 \rangle$ 7. Case: A is an end extension of B
  - $\langle 3 \rangle 1. \ a \leq_A x$

```
\langle 3 \rangle 2. \ a \leq x
```

 $\langle 2 \rangle 8$ . Case: B is an end extension of A

```
\langle 3 \rangle 1. Case: x \in A
```

- $\langle 4 \rangle 1. \ a \leq_A x$
- $\langle 4 \rangle 2$ .  $a \leq x$
- $\langle 3 \rangle 2$ . Case:  $x \in B A$ 
  - $\langle 4 \rangle 1. \ a \leq_B x$
  - $\langle 4 \rangle 2$ .  $a \leq x$
- $\langle 1 \rangle 6$ . For all  $A \in \mathcal{C}$ , W is an end extension of A.
  - $\langle 2 \rangle 1$ . For all  $x, y \in A$ , we have  $x \leq_A y$  if and only if  $x \leq_W y$ 
    - $\langle 3 \rangle 1$ . Let:  $x, y \in A$
    - $\langle 3 \rangle 2$ . If  $x \leq_A y$  then  $x \leq_W y$

PROOF: Immediate from definitions.

- $\langle 3 \rangle 3$ . If  $x \leq_W y$  then  $x \leq_A y$ 
  - $\langle 4 \rangle 1$ . Assume:  $x \leq_W y$
  - $\langle 4 \rangle 2$ . PICK  $B \in \mathcal{C}$  such that  $x \leq_B y$
  - $\langle 4 \rangle 3$ . Case: A is an end extension of B

PROOF: Then  $x \leq_A y$ .

 $\langle 4 \rangle 4$ . Case: B is an end extension of A

PROOF: Then  $x \leq_A y$ .

- $\langle 2 \rangle 2$ . For all  $x \in A$  and  $y \in W A$  we have x < y
  - $\langle 3 \rangle 1$ . Let:  $x \in A$  and  $y \in W A$
  - $\langle 3 \rangle 2$ . PICK  $B \in \mathcal{C}$  such that  $y \in B$
  - $\langle 3 \rangle 3$ . B is an end extension of A
  - $\langle 3 \rangle 4$ .  $x <_B y$
  - $\langle 3 \rangle 5$ .  $x <_W y$
- $\langle 1 \rangle 7$ . For all  $A \in \mathcal{C}$ , the ordinal of A is  $\leq$  the ordinal of W.

PROOF: Lemma 6.1.14.

- $\langle 1 \rangle 8$ . For any ordinal  $\alpha$ , if for all  $A \in \mathcal{C}$  the ordinal of A is  $\leq \alpha$ , then the ordinal of W is  $\leq \alpha$ .
  - $\langle 2 \rangle 1$ . Let:  $\alpha$  be an ordinal.
  - $\langle 2 \rangle 2$ . Assume: for all  $A \in \mathcal{C}$ , the ordinal of A is  $\leq \alpha$
  - $\langle 2 \rangle 3$ . Let:  $\beta$  be the ordinal of W
  - $\langle 2 \rangle 4$ . Let:  $E: W \approx \beta$  be the canonical isomorphism.
  - $\langle 2 \rangle$ 5. Assume: for a contradiction  $\alpha < \beta$
  - $\langle 2 \rangle$ 6. Let:  $a \in W$  be the element with  $E(a) = \alpha$
  - $\langle 2 \rangle 7$ . PICK  $A \in \mathcal{C}$  such that  $a \in A$
  - $\langle 2 \rangle 8$ . Let:  $\gamma$  be the ordinal of A and  $E_A: A \cong \gamma$  be the canonical isomorphism.
  - $\langle 2 \rangle 9$ . For all  $x \in A$  we have  $E_A(x) = E(x)$

PROOF: Transfinite induction on x.

- $\langle 2 \rangle 10. \ E_A(a) = \alpha$
- $\langle 2 \rangle 11. \ \alpha < \gamma$
- $\langle 2 \rangle 12$ . Q.E.D.

PROOF: This contradicts  $\langle 2 \rangle 2$ .

# 6.2 Successors

**Definition 6.2.1** (Successor). The *successor* of a set a is the set  $a^+ := a \cup \{a\}$ .

Proposition 6.2.2. A set a is a transitive set if and only if

$$\bigcup (a^+) = a .$$

Proof:

 $\langle 1 \rangle 1$ . If a is a transitive set then  $\bigcup (a^+) = a$ .

 $\langle 2 \rangle 1$ . Assume: a is a transitive set.

 $\langle 2 \rangle 2$ .  $\bigcup (a^+) \subseteq a$ 

 $\langle 3 \rangle 1$ . Let:  $x \in \bigcup (a^+)$ 

Prove:  $x \in a$ 

 $\langle 3 \rangle 2$ . PICK  $y \in a^+$  such that  $x \in y$ .

 $\langle 3 \rangle 3. \ y \in a \text{ or } y = a.$ 

 $\langle 3 \rangle 4$ . Case:  $y \in a$ 

PROOF: Then  $x \in a$  because a is a transitive set.

 $\langle 3 \rangle 5$ . Case: y = a

PROOF: Then  $x \in a$  immediately.

 $\langle 2 \rangle 3. \ a \subseteq \bigcup (a^+)$ 

PROOF: Since  $a \in a^+$ .

 $\langle 1 \rangle 2$ . If  $| J(a^+) = a$  then a is a transitive set.

 $\langle 2 \rangle 1$ . Assume:  $\bigcup (a^+) = a$ 

 $\langle 2 \rangle 2$ .  $\bigcup a \subseteq a$ 

Proof:

$$\bigcup a \subseteq \bigcup (a^{+})$$
 (Proposition 1.5.9)  
=  $a$  ( $\langle 2 \rangle 1$ )

 $\langle 2 \rangle 3$ . a is a transitive set.

Proof: Proposition 1.6.2.

**Proposition 6.2.3.** For any set a, we have a is a transitive set if and only if  $a^+$  is a transitive set.

PROOF:

 $\langle 1 \rangle 1$ . If a is a transitive set then  $a^+$  is a transitive set.

PROOF: If a is a transitive set then  $\bigcup (a^+) = a \subseteq a^+$  by Proposition 6.2.2 and so  $a^+$  is a transitive set.

 $\langle 1 \rangle 2$ . If  $a^+$  is a transitive set then a is a transitive set.

 $\langle 2 \rangle 1$ . Assume:  $a^+$  is a transitive set.

 $\langle 2 \rangle 2$ . Let:  $x \in y \in a$ 

 $\langle 2 \rangle 3. \ x \in y \in a^+$ 

 $\langle 2 \rangle 4. \ x \in a^+$ 

Proof:  $\langle 2 \rangle 1$ 

 $\langle 2 \rangle 5. \ x \neq a$ 

PROOF: From  $\langle 2 \rangle 2$  and the Axiom of Regularity.

$$\sqrt{2}6. \ x \in a$$

**Definition 6.2.4.** We write 0 for  $\emptyset$ , 1 for  $\emptyset^+$ , 2 for  $\emptyset^{++}$ , etc.

**Proposition 6.2.5.** For any set A we have  $\mathcal{P}A \approx 2^A$ .

PROOF: The function  $H: \mathcal{P}A \to 2^A$  defined by  $H(S)(a) = \{\emptyset\}$  if  $a \in S$  and  $\emptyset$  if  $a \notin S$  is a bijection.  $\square$ 

**Proposition 6.2.6.** For any ordinal number  $\alpha$  we have  $\alpha^+$  is an ordinal number.

#### Proof:

 $\langle 1 \rangle 1$ .  $\alpha^+$  is a transitive set.

PROOF: Proposition 6.2.3.

- $\langle 1 \rangle 2$ .  $\alpha^+$  is well-ordered by  $\in$ .
  - $\langle 2 \rangle 1$ . For all  $x, y, z \in \alpha^+$ , if  $x \in y \in z$  then  $x \in z$ 
    - $\langle 3 \rangle 1$ . Case:  $z = \alpha$

PROOF: Then  $x \in \alpha$  since  $\alpha$  is a transitive set.

 $\langle 3 \rangle 2$ . Case:  $z \in \alpha$ 

PROOF: Then  $x \in z$  since  $\alpha$  is well-ordered by  $\in$ .

- $\langle 2 \rangle 2$ . For all  $x, y \in \alpha^+$  we have  $x \in y$  or x = y or  $y \in x$ 
  - $\langle 3 \rangle 1$ . Case:  $x, y \in \alpha$

PROOF: The result follows because  $\alpha$  is well-ordered by  $\in$ .

 $\langle 3 \rangle 2$ . Case:  $x \in \alpha, y = \alpha$ 

PROOF: Then  $x \in y$ .

 $\langle 3 \rangle 3$ . Case:  $x = \alpha, y \in \alpha$ 

PROOF: Then  $y \in x$ .

 $\langle 3 \rangle 4$ . Case:  $x = \alpha, y = \alpha$ 

PROOF: Then x = y.

- $\langle 2 \rangle 3$ . Every nonempty subset of  $\alpha^+$  has an  $\in$ -least element.
  - $\langle 3 \rangle 1$ . Let:  $S \subseteq \alpha^+$  be nonempty
  - $\langle 3 \rangle 2$ . Case:  $S = \{\alpha\}$

PROOF:  $\alpha$  is least in S.

- $\langle 3 \rangle 3$ . Case:  $S \neq \{\alpha\}$ 
  - $\langle 4 \rangle 1$ .  $S \{\alpha\}$  is a nonempty subset of  $\alpha$
  - $\langle 4 \rangle 2$ . Let:  $\beta$  be least in  $S \{\alpha\}$
  - $\langle 4 \rangle 3$ .  $\beta$  is least in S.

**Proposition 6.2.7.** For ordinals  $\alpha$  and  $\beta$ , if  $\alpha^+ = \beta^+$  then  $\alpha = \beta$ .

PROOF: If 
$$\alpha^+ = \beta^+$$
 then 
$$\alpha = \bigcup (\alpha^+)$$
 (Proposition 6.2.2) 
$$= \bigcup (\beta^+)$$
 
$$= \beta$$
 (Proposition 6.2.2)

**Proposition 6.2.8.** For ordinals  $\alpha$  and  $\beta$ , we have  $\alpha < \beta$  if and only if  $\alpha^+ < \beta^+$ .

Proof:

$$\alpha < \beta \Leftrightarrow \alpha^+ \le \beta$$
$$\Leftrightarrow \alpha^+ < \beta^+$$

**Definition 6.2.9** (Successor Ordinal). An ordinal  $\alpha$  is a *successor ordinal* iff  $\alpha = \beta^+$  for some  $\beta$ .

**Definition 6.2.10** (Limit Ordinal). A *limit ordinal* is an ordinal that is neither 0 nor a successor ordinal.

**Proposition 6.2.11.** *If*  $\lambda$  *is a limit ordinal and*  $\beta < \lambda$  *then*  $\beta^+ < \lambda$ .

PROOF: Since  $\beta^+ \leq \lambda$  and  $\beta^+ \neq \lambda$ .  $\square$ 

# 6.3 The Well-Ordering Theorem and Zorn's Lemma

**Theorem 6.3.1** (Hartogs). For any set A, there exists an ordinal not dominated by A.

#### Proof:

- $\langle 1 \rangle$ 1. Let:  $\alpha$  be the class of all ordinals  $\beta$  such that  $\beta \leq A$  Prove:  $\alpha$  is a set.
- $\langle 1 \rangle 2$ . Let:  $W = \{(B, R) \mid B \subseteq A, R \text{ is a well ordering on } B\}$
- $\langle 1 \rangle 3$ .  $\alpha$  is the class of the ordinals of the elements of W.
  - $\langle 2 \rangle 1$ . For all  $(B,R) \in W$ , the ordinal of (B,R) is in  $\alpha$ .
    - $\langle 3 \rangle 1$ . Let:  $(B, R) \in W$
    - $\langle 3 \rangle 2$ . Let:  $\beta$  be the ordinal of (B, R)
    - $\langle 3 \rangle 3$ . Let:  $E: B \cong \beta$  be the canonical isomorphism.
    - $\langle 3 \rangle 4$ . Let:  $i: B \hookrightarrow A$  be the inclusion
    - $\langle 3 \rangle 5$ .  $i \circ E^{-1}$  is an injection  $\beta \to A$
    - $\langle 3 \rangle 6. \ \beta \in \alpha$
  - $\langle 2 \rangle 2$ . For all  $\beta \in \alpha$ , there exists  $(B,R) \in W$  such that  $\beta$  is the ordinal number of (B,R).
    - $\langle 3 \rangle 1$ . Let:  $\beta \in \alpha$
    - $\langle 3 \rangle 2$ . Pick an injection  $f: \beta \to A$
    - $\langle 3 \rangle 3$ . Define  $\leq$  on ran f by  $f(x) \leq f(y)$  iff  $x \leq y$
    - $\langle 3 \rangle 4$ .  $(\operatorname{ran} f, \leq) \in W$
  - $\langle 3 \rangle 5$ .  $\beta$  is the ordinal number of  $(\operatorname{ran} f, \leq)$
- $\langle 1 \rangle 4$ .  $\alpha$  is a set.

PROOF: By an Axiom of Replacement.

 $\langle 1 \rangle 5$ .  $\alpha$  is an ordinal.

PROOF: It is a transitive set of ordinals.

 $\langle 1 \rangle 6. \ \alpha \not \prec A$ 

PROOF: Since  $\alpha \notin \alpha$ .

П

**Theorem 6.3.2** (Numeration Theorem). Every set is equinumerous with some ordinal.

#### Proof:

- $\langle 1 \rangle 1$ . Let: A be any set.
- $\langle 1 \rangle 2$ . Pick an ordinal  $\alpha$  not dominated by A.
- $\langle 1 \rangle 3$ . Pick a choice function G for A.
- $\langle 1 \rangle 4$ . Pick  $e \notin A$
- $\langle 1 \rangle 5$ . Let:  $F : \alpha \to A \cup \{e\}$  by transfinite recursion:

$$F(\gamma) = \begin{cases} G(A - F(\{\delta \mid \delta < \gamma\}) & \text{if } A - F(\{\delta \mid \delta < \gamma\}) \neq \emptyset \\ e & \text{if } A - F(\{\delta \mid \delta < \gamma\}) = \emptyset \end{cases}$$

- $\langle 1 \rangle 6. \ e \in \operatorname{ran} F$ 
  - $\langle 2 \rangle 1$ . Assume: for a contradiction  $e \notin \operatorname{ran} F$
  - $\langle 2 \rangle 2$ . F is an injection  $\alpha \to A$ .
    - $\langle 3 \rangle$ 1. Let:  $\beta, \gamma \in \alpha$  with  $\beta \neq \gamma$ Prove:  $F(\beta) \neq F(\gamma)$
    - $\langle 3 \rangle$ 2. Assume: w.l.o.g.  $\beta < \gamma$
    - $\langle 3 \rangle 3. \ F(\gamma) \in A F(\{\delta \mid \delta < \gamma\})$
    - $\langle 3 \rangle 4. \ F(\gamma) \notin F(\{\delta \mid \delta < \gamma\})$
    - $\langle 3 \rangle 5. \ F(\gamma) \neq F(\beta)$
  - $\langle 2 \rangle 3$ . Q.E.D.

PROOF: This contradicts  $\langle 1 \rangle 2$ .

- $\langle 1 \rangle 7$ . Let:  $\delta$  be least such that  $F(\delta) = e$
- $\langle 1 \rangle 8. \ F \upharpoonright \delta : \delta \approx A$

**Theorem 6.3.3** (Well-Ordering Theorem). Any set can be well ordered.

#### Proof:

- (1)1. PICK an ordinal  $\delta$  and a bijection  $F: A \approx \delta$
- $\langle 1 \rangle 2$ . Define  $\leq$  on A by  $F(x) \leq F(y)$  iff  $x \leq y$  for  $x, y \in \delta$
- $\langle 1 \rangle 3. \leq \text{is a well ordering on } A.$

**Theorem 6.3.4** (Zorn's Lemma). Let  $\mathcal{A}$  be a set such that, for every chain  $\mathcal{B} \subseteq \mathcal{A}$ , we have  $\bigcup \mathcal{B} \in \mathcal{A}$ . Then  $\mathcal{A}$  has a maximal element.

#### Proof:

- $\langle 1 \rangle 1$ . Pick a well ordering  $\langle$  on  $\mathcal{A}$ .
- $\langle 1 \rangle 2$ . Let:  $F: \mathcal{A} \to 2$  be the function defined by transfinite recursion by:

$$F(A) = \begin{cases} 1 & \text{if } A \text{ includes every set } B < A \text{ for which } F(B) = 1 \\ 0 & \text{otherwise} \end{cases}$$

 $\langle 1 \rangle 3$ . Let:  $\mathcal{C} = \{ A \in \mathcal{A} \mid F(A) = 1 \}$ 

PROVE:  $\bigcup \mathcal{C}$  is a maximal element of  $\mathcal{A}$ 

- $\langle 1 \rangle 4$ . For all  $A \in \mathcal{A}$ , we have  $A \in \mathcal{C}$  iff  $\forall B < A.B \in \mathcal{C} \Rightarrow B \subseteq A$
- $\langle 1 \rangle 5$ . C is a chain.

```
 \begin{array}{l} \langle 2 \rangle 1. \text{ Let: } A, A' \in \mathcal{C} \\ \langle 2 \rangle 2. \text{ Assume: w.l.o.g. } A \leq A' \\ \langle 2 \rangle 3. A \subseteq A' \\ \text{ Proof: By } \langle 1 \rangle 4 \\ \langle 1 \rangle 6. \bigcup \mathcal{C} \in \mathcal{A} \\ \langle 1 \rangle 7. \bigcup \mathcal{C} \text{ is maximal in } \mathcal{A}. \\ \langle 2 \rangle 1. \text{ Let: } A \in \mathcal{A} \text{ and } \bigcup \mathcal{C} \subseteq A \\ \langle 2 \rangle 2. A \in \mathcal{C} \\ \text{ Proof: By } \langle 1 \rangle 4 \text{ since } \forall B \in \mathcal{C}.B \subseteq A. \\ \langle 2 \rangle 3. A \subseteq \bigcup \mathcal{C} \\ \langle 2 \rangle 4. A = \bigcup \mathcal{C} \\ \end{array}
```

**Proposition 6.3.5** (Teichmüller-Tukey Lemma). Let A be a nonempty set such that, for every B, we have  $B \in A$  if and only if every finite subset of B is a member of A. Then A has a maximal element.

#### Proof:

```
\langle 1 \rangle 1. For every chain \mathcal{B} \subseteq \mathcal{A}, we have \bigcup \mathcal{B} \in \mathcal{A} \langle 2 \rangle 1. Let: \mathcal{B} \subseteq \mathcal{A} be a chain. \langle 2 \rangle 2. Every finite subset of \bigcup \mathcal{B} is a member of \mathcal{A}. \langle 3 \rangle 1. Let: C be a finite subset of \bigcup \mathcal{B}. \langle 3 \rangle 2. Pick B \in \mathcal{B} such that C \subseteq B. \langle 3 \rangle 3. B \in \mathcal{A} \langle 3 \rangle 4. Every finite subset of B is in A. \langle 3 \rangle 5. C \in \mathcal{A} \langle 2 \rangle 3. \bigcup \mathcal{B} \in \mathcal{A}.
```

Proof: Zorn's lemma.  $\Box$ 

 $\langle 1 \rangle 2$ . Q.E.D.

**Theorem Schema 6.3.6.** For any class **A**, there exists a class **F** such that the following is a theorem:

If **A** is a proper class of ordinals, then  $\mathbf{F}: \mathbf{On} \to \mathbf{A}$  is an order isomorphism.

#### Proof:

- $\langle 1 \rangle 1$ . Define  $\mathbf{F} : \mathbf{On} \to \mathbf{A}$  by transfinite recursion as follows:  $\mathbf{F}(\alpha)$  is the least element of  $\mathbf{A}$  that is different from  $\mathbf{F}(\beta)$  for all  $\beta < \alpha$ .
- $\langle 1 \rangle 2$ . For all  $\alpha, \beta \in \mathbf{On}$ , if  $\alpha < \beta$  then  $\mathbf{F}(\alpha) < \mathbf{F}(\beta)$

PROOF: We have  $\mathbf{F}(\alpha) \neq \mathbf{F}(\beta)$  by the definition of  $\mathbf{F}(\beta)$ , and  $\mathbf{F}(\beta) \not< \mathbf{F}(\alpha)$  by the leastness of  $\mathbf{F}(\alpha)$ .

- $\langle 1 \rangle 3$ . **F** is surjective.
  - $\langle 2 \rangle 1$ . Let:  $\alpha \in \mathbf{A}$
  - $\langle 2 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall \beta \in \mathbf{A}$ , if  $\beta < \alpha$  then there exists  $\gamma$  such that  $\beta = \mathbf{F}(\gamma)$ .
  - $\langle 2 \rangle 3$ . Let:  $\gamma = \{ \delta \in \mathbf{On} \mid \mathbf{F}(\delta) < \alpha \}$
  - $\langle 2 \rangle 4$ .  $\gamma$  is a set.

PROOF: Axiom of Replacement applied to  $\alpha$ .

 $\langle 2 \rangle 5$ .  $\gamma$  is a transitive set.

PROOF: If  $\mathbf{F}(\delta) < \alpha$  and  $\epsilon < \delta$  then  $\mathbf{F}(\epsilon) < \alpha$  by  $\langle 1 \rangle 2$ .

 $\langle 2 \rangle 6$ .  $\gamma$  is an ordinal.

Proof: Proposition 6.1.8.

 $\langle 2 \rangle 7$ .  $\mathbf{F}(\gamma) = \alpha$ 

 $\langle 3 \rangle 1$ .  $\mathbf{F}(\gamma)$  is the least element of **A** different from  $\mathbf{F}(\delta)$  for all  $\delta < \gamma$ 

 $\langle 3 \rangle 2$ .  $\mathbf{F}(\gamma)$  is the least element of  $\mathbf{A}$  different from x for all  $x \in \mathbf{A}$  with  $x < \alpha$ 

 $\langle 3 \rangle 3. \ \mathbf{F}(\gamma) = \alpha$ 

# 6.4 Ordinal Operations

**Definition 6.4.1** (Ordinal Operation). An *ordinal operation* is a function  $\mathbf{On} \to \mathbf{On}$ .

**Definition 6.4.2** (Continuous). An ordinal operation  $\mathbf{T}: \mathbf{On} \to \mathbf{On}$  is *continuous* iff, for every limit ordinal  $\lambda$ , we have  $\mathbf{T}(\lambda) = \bigcup_{\alpha < \lambda} \mathbf{T}(\alpha)$ .

**Definition 6.4.3** (Normal). An ordinal operation is *normal* iff it is continuous and strictly monotone.

Proposition Schema 6.4.4. For any class T, the following is a theorem.

If **T** is a continuous ordinal operation and  $\forall \gamma. \mathbf{T}(\gamma) < \mathbf{T}(\gamma^+)$ , then **T** is normal.

Proof:

 $\langle 1 \rangle 1$ . Let:  $P[\beta]$  be the property  $\forall \gamma < \beta. \mathbf{T}(\gamma) < \mathbf{T}(\beta)$ 

 $\langle 1 \rangle 2$ . P[0]

Proof: Vacuous.

 $\langle 1 \rangle 3$ . For any ordinal  $\gamma$ , if  $P[\gamma]$  then  $P[\gamma^+]$ 

 $\langle 2 \rangle 1$ . Assume:  $P[\gamma]$ 

 $\langle 2 \rangle 2$ . Let:  $\delta < \gamma^+$ 

 $\langle 2 \rangle 3$ . Case:  $\delta < \gamma$ 

PROOF: Then  $\mathbf{T}(\delta) < \mathbf{T}(\gamma) < \mathbf{T}(\gamma^+)$ .

 $\langle 2 \rangle 4$ . Case:  $\delta = \gamma$ 

PROOF: Then  $\mathbf{T}(\delta) = \mathbf{T}(\gamma) < \mathbf{T}(\gamma^+)$ .

 $\langle 1 \rangle 4$ . For any limit ordinal  $\lambda$ , if  $\forall \gamma < \lambda . P[\gamma]$  then  $P[\lambda]$ .

 $\langle 2 \rangle 1$ . Assume:  $\forall \gamma < \lambda . P[\gamma]$ 

 $\langle 2 \rangle 2$ . Let:  $\delta < \lambda$ 

 $\langle 2 \rangle 3. \ \mathbf{T}(\delta) < \mathbf{T}(\lambda)$ 

$$\mathbf{T}(\delta) < \mathbf{T}(\delta^{+})$$

$$\leq \bigcup_{\epsilon < \lambda} \mathbf{T}(\epsilon)$$

$$= \mathbf{T}(\lambda)$$

**Proposition Schema 6.4.5.** For any class **T**, the following is a theorem:

Assume **T** is a normal ordinal operation. For every ordinal  $\alpha$ , we have  $\alpha \leq \mathbf{T}(\alpha)$ .

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $\gamma$  be an ordinal.
- $\langle 1 \rangle 2$ . Assume: as induction hypothesis  $\forall \delta < \gamma. \mathbf{T}(\delta) \geq \delta$
- $\langle 1 \rangle 3$ . For all  $\delta < \gamma$  we have  $\delta < \mathbf{T}(\gamma)$

PROOF: **T** is strictly monotone.

 $\langle 1 \rangle 4. \ \gamma \leq \mathbf{T}(\gamma)$ 

**Proposition Schema 6.4.6.** For any class **T**, the following is a theorem:

Assume **T** is a normal ordinal operation. For any ordinal  $\beta \geq \mathbf{T}(0)$ , there exists a greatest ordinal  $\gamma$  such that  $\mathbf{T}(\gamma) \leq \beta$ .

#### Proof:

- $\langle 1 \rangle 1$ . There exists  $\gamma$  such that  $\mathbf{T}(\gamma) > \beta$ 
  - $\langle 2 \rangle 1$ . For all  $\gamma$  we have  $\mathbf{T}(\gamma) \geq \gamma$

Proof: Proposition 6.4.5.

- $\langle 2 \rangle 2$ .  $\mathbf{T}(\beta^+) > \beta$
- $\langle 1 \rangle 2$ . Let:  $\delta$  be least such that  $\mathbf{T}(\delta) > \beta$
- $\langle 1 \rangle 3$ .  $\delta$  is a successor ordinal.
  - $\langle 2 \rangle 1. \ \delta \neq 0$

PROOF: Since  $\mathbf{T}(0) \leq \beta$ .

- $\langle 2 \rangle 2$ .  $\delta$  is not a limit ordinal.
  - $\langle 3 \rangle 1$ . Assume: for a contradiction  $\delta$  is a limit ordinal.
  - $\langle 3 \rangle 2. \ \beta < \bigcup_{\epsilon < \delta} \mathbf{T}(\epsilon)$

PROOF: T is continuous.

- $\langle 3 \rangle 3$ . There exists  $\epsilon < \delta$  such that  $\beta < \mathbf{T}(\epsilon)$
- $\langle 3 \rangle 4$ . Q.E.D.

PROOF: This contradicts the minimality of  $\delta$ .

- $\langle 1 \rangle 4$ . Let:  $\delta = \gamma^+$
- $\langle 1 \rangle$ 5.  $\gamma$  is greatest such that  $\mathbf{T}(\gamma) \leq \beta$

**Theorem Schema 6.4.7.** For any class **T**, the following is a theorem:

Assume that T is a normal ordinal operation. For any nonempty set of ordinals S, we have

$$\mathbf{T}(\sup S) = \sup_{\alpha \in S} \mathbf{T}(\alpha) .$$

#### Proof:

 $\langle 1 \rangle 1. \ \forall \alpha \in S. \mathbf{T}(\alpha) \leq \mathbf{T}(\sup S)$ 

PROOF: Since T is monotone.

 $\langle 1 \rangle 2$ . For any ordinal  $\beta$ , if  $\forall \alpha \in S.\mathbf{T}(\alpha) \leq \beta$ , then  $\mathbf{T}(\sup S) \leq \beta$ 

```
\langle 2 \rangle 1. Let: \beta be an ordinal.
    \langle 2 \rangle 2. Let: \gamma = \sup S
    \langle 2 \rangle 3. Assume: \forall \alpha \in S.\mathbf{T}(\alpha) \leq \beta
    \langle 2 \rangle 4. Case: \gamma is 0 or a successor ordinal
        PROOF: Then we must have \gamma \in S so \mathbf{T}(\gamma) \leq \beta from \langle 2 \rangle 3.
    \langle 2 \rangle5. Case: \gamma is a limit ordinal
        \langle 3 \rangle 1. \mathbf{T}(\gamma) = \sup_{\alpha < \gamma} \mathbf{T}(\alpha)
           PROOF: T is continuous.
        \langle 3 \rangle 2. Assume: for a contradiction \beta < \mathbf{T}(\gamma)
        \langle 3 \rangle 3. Pick \alpha < \gamma such that \beta < \mathbf{T}(\alpha)
           Proof: \langle 3 \rangle 1, \langle 3 \rangle 2
        \langle 3 \rangle 4. Pick \alpha' \in S such that \alpha < \alpha'
           Proof: \langle 2 \rangle 2, \langle 3 \rangle 3
        \langle 3 \rangle 5. \ \beta < \mathbf{T}(\alpha') \leq \beta
           PROOF: T is strictly monotone, \langle 3 \rangle 3, \langle 3 \rangle 4, \langle 2 \rangle 3.
        \langle 3 \rangle 6. Q.E.D.
           Proof: This is a contradiction.
П
Proposition 6.4.8. For any classes A and T, the following is a theorem:
      Assume A is a proper class of ordinals such that, for every set S \subseteq \mathbf{A}, we
have | \ | \ S \in \mathbf{A}. Assume T is the unique order isomorphism \mathbf{On} \cong \mathbf{A}. Then T
is normal.
Proof:
\langle 1 \rangle 1. T is strictly monotone.
    PROOF: Since it is an order isomorphism.
\langle 1 \rangle 2. T is continuous.
    \langle 2 \rangle 1. Let: \lambda be a limit ordinal.
    \langle 2 \rangle 2. \mathbf{T}'(\lambda) is the least member of A that is greater than \mathbf{T}'(\alpha) for all \alpha < \lambda
    \langle 2 \rangle 3. \ \mathbf{T}'(\lambda) = \sup_{\alpha < \lambda} \mathbf{T}'(\alpha)
Proposition Schema 6.4.9. For any class T, the following is a theorem:
      If T is a normal ordinal operation, then for any limit ordinal \lambda, we have
\mathbf{T}(\lambda) is a limit ordinal.
Proof:
\langle 1 \rangle 1. \mathbf{T}(\lambda) \neq 0
    PROOF: Since 0 \leq \mathbf{T}(0) < \mathbf{T}(\lambda).
\langle 1 \rangle 2. \mathbf{T}(\lambda) is not a successor ordinal.
    \langle 2 \rangle 1. Assume: for a contradiction \mathbf{T}(\lambda) = \alpha^+
    \langle 2 \rangle 2. \alpha < \mathbf{T}(\lambda) = \sup_{\beta < \lambda} \mathbf{T}(\beta)
    \langle 2 \rangle 3. PICK \beta < \lambda such that \alpha < \mathbf{T}(\beta)
    \langle 2 \rangle 4. \ \alpha^+ \leq \mathbf{T}(\beta) < \mathbf{T}(\lambda)
    \langle 2 \rangle 5. Q.E.D.
       PROOF: This is a contradiction.
```

### 6.5 Ordinal Arithmetic

### 6.5.1 Addition

**Definition 6.5.1.** Let A and B be disjoint well-ordered sets. The *concatenation* of A and B is the set  $A \cup B$  under the relation:

- if  $a, a' \in A$  then  $a \leq a'$  iff  $a \leq a'$  in A
- if  $b, b' \in B$  then  $b \le b'$  iff  $b \le b'$  in B
- if  $a \in A$  and  $b \in B$  then  $a \le b$  and  $b \not \le a$ .

**Proposition 6.5.2.** If A and B are disjoint well-ordered sets, then their concatenation is well-ordered.

#### Proof:

 $\langle 1 \rangle 1. \leq \text{is reflexive.}$ 

PROOF: For all  $a \in A$  we have  $a \le a$ , and for all  $b \in B$  we have  $b \le b$ .

- $\langle 1 \rangle 2$ . < is antisymmetric.
  - $\langle 2 \rangle 1$ . Assume:  $x \leq y \leq x$
  - $\langle 2 \rangle 2$ . Case:  $x, y \in A$

PROOF: Then x = y since the order on A is antisymmetric.

 $\langle 2 \rangle 3$ . Case:  $x \in A$  and  $y \in B$ 

PROOF: This is impossible as it would imply  $y \not\leq x$ .

 $\langle 2 \rangle 4$ . Case:  $x \in B$  and  $y \in A$ 

PROOF: This is impossible as it would imply  $x \not\leq y$ .

 $\langle 2 \rangle$ 5. Case:  $x, y \in B$ 

PROOF: Then x = y since the order on B is antisymmetric.

- $\langle 1 \rangle 3. \leq \text{is transitive.}$ 
  - $\langle 2 \rangle 1$ . Assume:  $x \leq y \leq z$
  - $\langle 2 \rangle 2$ . Case:  $x, z \in A$

PROOF: In this case  $y \in A$  since  $y \le z$ , and so  $x \le z$  since the order on A is transitive.

 $\langle 2 \rangle 3$ . Case:  $x \in A$  and  $z \in B$ 

PROOF: Then  $x \leq z$  immediately.

 $\langle 2 \rangle 4$ . Case:  $x \in B$  and  $z \in A$ 

PROOF: This is impossible because we have  $y \notin A$  since  $x \leq y$  and  $y \notin B$  since  $y \leq z$ .

 $\langle 2 \rangle$ 5. Case:  $x, z \in B$ 

PROOF: In this case  $y \in B$  since  $x \le y$ , and so  $x \le z$  since the order on B is transitive.

- $\langle 1 \rangle 4$ .  $\leq$  is total.
  - $\langle 2 \rangle 1$ . Let:  $x, y \in A \cup B$
  - $\langle 2 \rangle 2$ . Case:  $x, y \in A$

PROOF: Then  $x \leq y$  or  $y \leq x$  because the order on A is total.

 $\langle 2 \rangle 3$ . Case:  $x \in A$  and  $y \in B$ 

PROOF: Then  $x \leq y$ .

 $\langle 2 \rangle 4$ . Case:  $x \in B$  and  $y \in A$ 

PROOF: Then  $y \le x$ .  $\langle 2 \rangle 5$ . Case:  $x, y \in B$ 

PROOF: Then  $x \leq y$  or  $y \leq x$  because the order on B is total.

- $\langle 1 \rangle 5$ . Every nonempty subset of  $A \cup B$  has a least element.
  - $\langle 2 \rangle 1$ . Let: S be a nonempty subset of  $A \cup B$
  - $\langle 2 \rangle 2$ . Case:  $S \cap A = \emptyset$

PROOF: Then  $S \subseteq B$  and so S has a least element.

 $\langle 2 \rangle 3$ . Case:  $S \cap A \neq \emptyset$ 

PROOF: The least element of  $S \cap A$  is the least element of S.

**Definition 6.5.3** (Ordinal Addition). Let  $\alpha$  and  $\beta$  be ordinal numbers. Then  $\alpha + \beta$  is the ordinal number of the concatenation of A and B, where A is any well ordered set with ordinal  $\alpha$  and B is any well ordered set with ordinal  $\beta$ .

**Theorem 6.5.4** (Associative Law for Addition). For any ordinals  $\rho$ ,  $\sigma$  and  $\tau$ , we have

$$\rho + (\sigma + \tau) = (\rho + \sigma) + \tau .$$

PROOF: Given disjoint well ordered sets A, B and C, the concatenation of A with (the concatenation of B and C) is the same as the concatenation of (the concatenation of A and B) and C.  $\square$ 

**Theorem 6.5.5.** For any ordinal  $\rho$  we have

$$\rho + 0 = 0 + \rho = \rho$$
.

PROOF: For any well ordered set A, the concatenation of A with  $\emptyset$  is A, and the concatenation of  $\emptyset$  with A is A.  $\square$ 

**Theorem 6.5.6.** For any ordinal  $\alpha$  we have  $\alpha + 1 = \alpha^+$ .

PROOF: Since  $\alpha^+$  is the concatenation of  $\alpha$  and  $\{\alpha\}$ .  $\square$ 

**Theorem 6.5.7.** For any ordinal  $\alpha$ , the operation that maps  $\beta$  to  $\alpha + \beta$  is normal.

PROOF

- $\langle 1 \rangle 1$ . For any limit ordinal  $\lambda$ , we have  $\alpha + \lambda = \sup_{\beta < \lambda} (\alpha + \beta)$ .
  - $\langle 2 \rangle 1$ . Let:  $\lambda$  be a limit ordinal.
  - $\langle 2 \rangle 2$ .  $(\{0\} \times \alpha) \cup (\{1\} \times \lambda) = \bigcup_{\beta \in \lambda} ((\{0\} \times \alpha) \cup (\{1\} \times \beta))$ , where the order on the right hand side is as in Lemma 6.1.15.

$$(\{0\} \times \alpha) \cup (\{1\} \times \lambda) = (\{0\} \times \alpha) \cup (\{1\} \times \bigcup_{\beta < \lambda} \beta)$$
$$= (\{0\} \times \alpha) \cup \bigcup_{\beta < \lambda} (\{1\} \times \beta)$$
$$= \bigcup_{\beta < \lambda} ((\{0\} \times \alpha) \cup (\{1\} \times \beta))$$

 $\langle 1 \rangle$ 2. For any ordinal  $\beta$  we have  $\alpha + \beta < \alpha + \beta^+$ PROOF: Since  $\alpha + \beta^+ = \alpha + \beta + 1 = (\alpha + \beta)^+$ .

**Corollary 6.5.7.1.** For any ordinals  $\alpha$ ,  $\beta$ ,  $\gamma$ , we have  $\beta < \gamma$  if and only if  $\alpha + \beta < \alpha + \gamma$ .

**Corollary 6.5.7.2** (Left Cancellation for Addition). For any ordinals  $\alpha$ ,  $\beta$  and  $\gamma$ , if  $\alpha + \beta = \alpha + \gamma$  then  $\beta = \gamma$ .

**Theorem 6.5.8.** For any ordinals  $\alpha$ ,  $\beta$ ,  $\gamma$ , if  $\beta \leq \gamma$  then  $\beta + \alpha \leq \gamma + \alpha$ .

PROOF: Transfinite induction on  $\alpha$ .

**Theorem 6.5.9** (Subtraction Theorem). Let  $\alpha$  and  $\beta$  be ordinals with  $\alpha \leq \beta$ . Then there exists a unique ordinal  $\delta$  such that  $\alpha + \delta = \beta$ .

Proof:

- $\langle 1 \rangle 1$ . For all ordinals  $\alpha$  and  $\beta$  with  $\alpha \leq \beta$ , there exists  $\delta$  such that  $\alpha + \delta = \beta$ 
  - $\langle 2 \rangle 1$ . Let:  $\alpha$  and  $\beta$  be ordinals with  $\alpha \leq \beta$
  - $\langle 2 \rangle 2$ . Let:  $\delta$  be the greatest ordinal such that  $\alpha + \delta \leq \beta$

Proof: Proposition 6.4.6.

 $\langle 2 \rangle 3. \ \alpha + \delta = \beta$ 

PROOF: If  $\alpha + \delta < \beta$  then  $\alpha + \delta + 1 \le \beta$  contradicting the greatestness of  $\delta$ .  $\langle 1 \rangle 2$ . Q.E.D.

PROOF: Uniqueness follows from the Left Cancellation Law.  $\Box$ 

6.5.2 Multiplication

**Definition 6.5.10** (Ordinal Multiplication). Let  $\alpha$  and  $\beta$  be ordinal numbers. Then  $\alpha\beta$  is the ordinal number of  $A \times B$  under the lexicographic order, where A is any well ordered set with ordinal  $\alpha$  and B is any well ordered set with ordinal  $\beta$ .

This is well defined by Proposition 5.3.3.

**Theorem 6.5.11** (Associative Law). For any ordinals  $\rho$ ,  $\sigma$  and  $\tau$ , we have

$$\rho(\sigma\tau) = (\rho\sigma)\tau .$$

PROOF: Let A, B and C be well ordered sets with ordinals  $\rho$ ,  $\sigma$  and  $\tau$ . Then both  $\rho(\sigma\tau)$  and  $(\rho\sigma)\tau$  are the ordinal of  $A\times B\times C$  under  $(a,b,c)\leq (a',b',c')\Leftrightarrow a\leq a'\vee (a=a'\wedge b\leq b')\vee (a=a'\wedge b=b'\wedge c\leq c')$ .

**Theorem 6.5.12** (Left Distributive Law). For any ordinals  $\rho$ ,  $\sigma$  and  $\tau$ , we have

$$\rho(\sigma + \tau) = \rho\sigma + \rho\tau$$

PROOF: Let A, B and C be well ordered sets with ordinals  $\rho$ ,  $\sigma$  and  $\tau$  and with  $B \cap C = \emptyset$ . Then both  $\rho(\sigma + \tau)$  and  $\rho\sigma + \rho\tau$  are the ordinal of  $A \times (B \cup C)$  under the lexicographic ordering.  $\square$ 

**Theorem 6.5.13.** For any ordinal  $\rho$  we have  $\rho 0 = 0\rho = 0$ .

PROOF: For any well ordered set A we have  $A \times \emptyset = \emptyset \times A = \emptyset$ .  $\square$ 

**Theorem 6.5.14.** For any ordinal  $\rho$  we have  $\rho 1 = 1\rho = \rho$ .

Proof: Easy.  $\square$ 

**Theorem 6.5.15.** For any ordinals  $\rho$  and  $\sigma$ , if  $\rho\sigma = 0$  then  $\rho = 0$  or  $\sigma = 0$ .

PROOF: If  $A \times B = \emptyset$  then  $A = \emptyset$  or  $B = \emptyset$ .  $\square$ 

**Theorem 6.5.16.** For any non-zero ordinal  $\alpha$ , the operation that maps  $\beta$  to  $\alpha\beta$  is normal.

#### Proof:

- $\langle 1 \rangle 1$ . For any limit ordinal  $\lambda$ , we have  $\alpha \lambda = \bigcup_{\beta < \lambda} \alpha \beta$ 
  - $\langle 2 \rangle$ 1. Let:  $\lambda$  be a limit ordinal
  - $\langle 2 \rangle 2$ .  $\alpha \times \lambda = \bigcup_{\beta < \lambda} (\alpha \times \beta)$  as well-ordered sets
- $\langle 1 \rangle 2$ . For any ordinal  $\beta$  we have  $\alpha \beta < \alpha \beta^+$

PROOF:  $\alpha \beta^+ = \alpha \beta + \alpha > \alpha \beta$ 

П

**Corollary 6.5.16.1.** For any ordinals  $\alpha$ ,  $\beta$ ,  $\gamma$ , if  $\alpha \neq 0$  then  $\beta < \gamma$  if and only if  $\alpha\beta < \alpha\gamma$ .

**Corollary 6.5.16.2** (Left Cancellation for Multiplication). For any ordinals  $\alpha$ ,  $\beta$ ,  $\gamma$ , if  $\alpha \neq 0$  and  $\alpha\beta = \alpha\gamma$  then  $\beta = \gamma$ .

**Theorem 6.5.17.** For any ordinals  $\alpha$ ,  $\beta$  and  $\gamma$ , if  $\beta \leq \gamma$  then  $\beta \alpha \leq \gamma \alpha$ .

PROOF: Transfinite induction on  $\alpha$ .

**Theorem 6.5.18** (Division Theorem). Let  $\alpha$  and  $\delta$  be ordinal numbers with  $\delta \neq 0$ . Then there exist unique ordinals  $\beta$  and  $\gamma$  with  $\gamma < \delta$  and

$$\alpha = \delta \beta + \gamma$$
.

- $\langle 1 \rangle 1$ . For any ordinal numbers  $\alpha$  and  $\delta$  with  $\delta \neq 0$ , there exist ordinals  $\beta$  and  $\gamma$  such that  $\gamma < \delta$  and  $\alpha = \delta \beta + \gamma$ 
  - $\langle 2 \rangle 1$ . Let:  $\alpha$  and  $\delta$  be ordinals with  $\delta \neq 0$
  - $\langle 2 \rangle$ 2. Let:  $\beta$  be the greatest ordinal such that  $\delta \beta \leq \alpha$  Proof: Proposition 6.4.6.
  - $\langle 2 \rangle$ 3. There exists an ordinal  $\gamma$  such that  $\alpha = \delta \beta + \gamma$ PROOF: Subtraction Theorem

$$\langle 1 \rangle 2$$
. For any ordinals  $\delta$ ,  $\beta$ ,  $\beta'$ ,  $\gamma$ ,  $\gamma'$ , if  $\delta \beta + \gamma = \delta \beta' + \gamma'$  and  $\delta \neq 0$  and  $\gamma$ ,  $\gamma' < \delta$  then  $\beta = \beta'$  and  $\gamma = \gamma'$ 

$$\langle 2 \rangle 1$$
. Let:  $\delta$ ,  $\beta$ ,  $\beta'$ ,  $\gamma$ ,  $\gamma'$  be ordinals.

$$\langle 2 \rangle 2$$
. Assume:  $\delta \neq 0$  and  $\delta \beta + \gamma = \delta \beta' + \gamma'$ 

$$\langle 2 \rangle 3. \ \beta = \beta'$$

$$\langle 3 \rangle 1. \ \beta \not< \beta'$$

PROOF: If  $\beta < \beta'$  then

$$\delta\beta' + \gamma' \ge \delta\beta'$$

$$\ge \delta(\beta + 1)$$

$$= \delta\beta + \delta$$

$$> \delta\beta + \gamma$$

 $\langle 3 \rangle 2. \ \beta' \not < \beta$ 

PROOF: Similar.

 $\langle 2 \rangle 4. \ \gamma = \gamma'$ 

PROOF: By Cancellation.

#### 6.5.3Exponentiation

**Definition 6.5.19.** Given ordinals  $\alpha$  and  $\beta$ , define the ordinal  $\alpha^{\beta}$  as follows:

$$0^{\alpha} := 0 \qquad (\alpha > 0)$$

$$\alpha^{0} := 1$$

$$\alpha^{\beta^{+}} := \alpha^{\beta} \alpha \qquad (\alpha > 0)$$

$$\alpha^{\lambda} := \sup_{\beta < \lambda} \alpha^{\beta} \qquad (\alpha > 0, \lambda \text{ a limit ordinal})$$

**Theorem 6.5.20.** Let  $\alpha$  be an ordinal  $\geq 2$ . The operation that maps  $\beta$  to  $\alpha^{\beta}$ is normal.

Proof:

 $\langle 1 \rangle 1$ . For  $\lambda$  a limit ordinal we have  $\alpha^{\lambda} = \sup_{\beta < \lambda} \alpha^{\beta}$ 

PROOF: By definition.

 $\langle 1 \rangle 2$ . For any ordinal  $\beta$  we have  $\alpha^{\beta} < \alpha^{\beta^+}$ 

PROOF: We have  $\alpha^{\beta^+} = \alpha^{\beta} \alpha > \alpha^{\beta}$  by Theorem 6.5.16 since  $\alpha > 1$  and  $\alpha^{\beta} \neq 0$ .

**Corollary 6.5.20.1.** For any ordinals  $\alpha$ ,  $\beta$ ,  $\gamma$ , if  $\alpha \geq 2$  then  $\beta < \gamma$  if and only if  $\alpha^{\beta} < \alpha^{\gamma}$ .

Corollary 6.5.20.2 (Cancellation for Exponentiation). For any ordinals  $\alpha$ ,  $\beta$ ,  $\gamma$ , if  $\alpha \geq 2$  and  $\alpha^{\beta} = \alpha^{\gamma}$  then  $\beta = \gamma$ .

**Theorem 6.5.21.** For any ordinals  $\alpha$ ,  $\beta$  and  $\gamma$ , if  $\beta \leq \gamma$  then  $\beta^{\alpha} \leq \gamma^{\alpha}$ .

PROOF: Transfinite induction on  $\alpha$ .

**Theorem 6.5.22** (Logarithm Theorem). Let  $\alpha$  and  $\beta$  be ordinal numbers with  $\alpha \neq 0$  and  $\beta > 1$ . Then there exist unique ordinals  $\gamma$ ,  $\delta$  and  $\rho$  such that

$$\alpha = \beta^{\gamma} \delta + \rho, \qquad 0 \neq \delta < \beta, \qquad \rho < \beta^{\gamma}.$$

Proof:

(1)1. For any ordinals  $\alpha$  and  $\beta$  with  $\alpha \neq 0$  and  $\beta > 1$ , there exist ordinals  $\gamma$ ,  $\delta$ ,  $\rho$  such that

$$\alpha = \beta^{\gamma} \delta + \rho, \qquad 0 \neq \delta < \beta, \qquad \rho < \beta^{\gamma}.$$

- $\langle 2 \rangle 1$ . Let:  $\alpha$  and  $\beta$  be ordinals with  $\alpha \neq 0$  and  $\beta > 1$ .
- $\langle 2 \rangle 2$ . Let:  $\gamma$  be the greatest ordinal such that  $\beta^{\gamma} \leq \alpha$ . Proof: Proposition 6.4.6.
- $\langle 2 \rangle 3$ . Let:  $\delta$  and  $\rho$  be the unique ordinals with  $\rho < \beta^{\gamma}$  such that  $\alpha = \beta^{\gamma} \delta + \rho$ . PROOF: By the Division Theorem.
- $\langle 2 \rangle 4. \ \delta \neq 0$

PROOF: If  $\delta = 0$  then  $\alpha = \beta^{\gamma}0 + \rho = \rho < \beta^{\gamma} \le \alpha$  which is a contradiction.  $\langle 2 \rangle 5. \ \delta < \beta$ 

PROOF: If  $\beta \leq \delta$  then  $\alpha \geq \beta^{\gamma} \delta \geq \beta^{\gamma} \beta = \beta^{\gamma+1}$ , contradicting the greatest-

- $\langle 1 \rangle 2$ . If  $\beta^{\gamma} \delta + \rho = \beta^{\gamma'} \delta' + \rho'$  with  $\beta > 1$ ,  $0 \neq \delta < \beta$ ,  $0 \neq \delta' < \beta$ ,  $\rho < \beta^{\gamma}$  and  $\rho' < \beta^{\gamma'}$ , then  $\gamma = \gamma'$ ,  $\delta = \delta'$  and  $\rho = \rho'$ .
  - $\langle 2 \rangle 1$ . Let:  $\alpha = \beta^{\gamma} \delta + \rho = \beta^{\gamma'} \delta' + \rho'$

  - $\begin{array}{l} \langle 2 \rangle 2. \ \beta^{\gamma} \leq \alpha < \beta^{\gamma+1} \\ \langle 2 \rangle 3. \ \beta^{\gamma'} \leq \alpha < \beta^{\gamma'+1} \\ \langle 2 \rangle 4. \ \beta^{\gamma} < \beta^{\gamma'+1} \ \text{and} \ \beta^{\gamma'} < \beta^{\gamma+1} \end{array}$
  - $\langle 2 \rangle 5$ .  $\gamma < \gamma' + 1$  and  $\gamma' < \gamma + 1$
  - $\langle 2 \rangle 6. \ \gamma = \gamma'$
  - $\langle 2 \rangle 7$ .  $\delta = \delta'$  and  $\rho = \rho'$

PROOF: By the Division Theorem.

**Theorem 6.5.23.** For any ordinal numbers  $\alpha$ ,  $\beta$ ,  $\gamma$ , we have

$$\alpha^{\beta+\gamma}=\alpha^{\beta}\alpha^{\gamma}$$
 .

Proof:

- (1)1. Let:  $P[\gamma]$  be the property: for any ordinals  $\alpha$  and  $\beta$  we have  $\alpha^{\beta+\gamma} = \alpha^{\beta}\alpha^{\gamma}$  $\langle 1 \rangle 2$ . P[0]
  - Proof:

$$\alpha^{\beta+0} = \alpha^{\beta}$$
$$= \alpha^{\beta} 1$$
$$= \alpha^{\beta} \alpha^{0}$$

 $\langle 1 \rangle 3$ . For all  $\gamma$ , if  $P[\gamma]$  then  $P[\gamma + 1]$ 

$$\begin{split} \alpha^{\beta+\gamma+1} &= \alpha^{\beta+\gamma}\alpha \\ &= \alpha^{\beta}\alpha^{\gamma}\alpha \\ &= \alpha^{\beta}\alpha^{\gamma+1} \end{split} \tag{induction hypothesis}$$

#### 6.5. ORDINAL ARITHMETIC

- $\langle 1 \rangle 4$ . For any limit ordinal  $\lambda$ , if  $\forall \gamma < \lambda . P[\gamma]$  then  $P[\lambda]$ .
  - $\langle 2 \rangle$ 1. Let:  $\lambda$  be a limit ordinal
  - $\langle 2 \rangle 2$ . Assume:  $\forall \gamma < \lambda . P[\gamma]$
  - $\langle 2 \rangle 3$ . Let:  $\alpha$  and  $\beta$  be any ordinals.
  - $\langle 2 \rangle 4$ . Case:  $\alpha = 0$

PROOF: We have  $\alpha^{\beta+\lambda} = \alpha^{\beta}\alpha^{\lambda} = 0$ .

 $\langle 2 \rangle$ 5. Case:  $\alpha = 1$ 

PROOF: We have  $\alpha^{\beta+\lambda} = \alpha^{\beta}\alpha^{\lambda} = 1$ .

 $\langle 2 \rangle 6$ . Case:  $\alpha > 1$ 

Proof:

$$\begin{split} \alpha^{\beta+\lambda} &= \alpha^{\sup_{\gamma < \lambda} (\beta+\gamma)} \\ &= \sup_{\gamma < \lambda} \alpha^{\beta+\gamma} & \text{(Theorem 6.4.7)} \\ &= \sup_{\gamma < \lambda} \alpha^{\beta} \alpha^{\gamma} & \text{($\langle 2 \rangle 2$)} \\ &= \alpha^{\beta} \sup_{\gamma < \lambda} \alpha^{\gamma} & \text{(Theorem 6.4.7)} \\ &= \alpha^{\beta} \alpha^{\lambda} \end{split}$$

**Theorem 6.5.24.** For any ordinal numbers  $\alpha$ ,  $\beta$  and  $\gamma$ , we have

$$(\alpha^{\beta})^{\gamma} = \alpha^{\beta\gamma} .$$

Proof:

- (1)1. Let:  $P[\gamma]$  be the property: For any ordinals  $\alpha$  and  $\beta$ , we have  $(\alpha^{\beta})^{\gamma}=\alpha^{\beta\gamma}$
- $\langle 1 \rangle 2$ . P[0] Proof:

$$(\alpha^{\beta})^0 = 1$$
$$= \alpha^{\beta 0}$$

 $\langle 1 \rangle 3. \ \forall \gamma \in \mathbf{On}.P[\gamma] \Rightarrow P[\gamma + 1]$ 

$$(\alpha^{\beta})^{\gamma+1} = (\alpha^{\beta})^{\gamma} \alpha^{\beta}$$
$$= \alpha^{\beta\gamma} \alpha^{\beta}$$
$$= \alpha^{\beta\gamma+\beta}$$
$$= \alpha^{\beta(\gamma+1)}$$

- $\langle 1 \rangle 4$ . For any limit ordinal  $\lambda$ , if  $\forall \gamma < \lambda . P[\gamma]$  then  $P[\lambda]$ .
  - $\langle 2 \rangle 1$ . Let:  $\lambda$  be a limit ordinal.
  - $\langle 2 \rangle 2$ . Assume:  $\forall \gamma < \lambda . P[\gamma]$
  - $\langle 2 \rangle 3$ . Let:  $\alpha$  and  $\beta$  be any ordinals.
  - $\langle 2 \rangle 4$ . Case:  $\alpha = 0$  and  $\beta = 0$

Proof:

$$(0^{\beta})^{\lambda} = 1^{\lambda}$$

$$= 1$$

$$= 0^{0}$$

$$= 0^{0\lambda}$$

$$\langle 2 \rangle$$
5. Case:  $\alpha = 0$  and  $\beta \neq 0$   
Proof:  $(0^{\beta})^{\lambda} = 0^{\beta\lambda} = 0$ .  
 $\langle 2 \rangle$ 6. Case:  $\alpha = 1$   
Proof:  $(1^{\beta})^{\lambda} = 1^{\beta\lambda} = 1$ 

$$\langle 2 \rangle 6$$
. Case:  $\alpha = 1$ 

PROOF: 
$$(1^{\beta})^{\lambda} = 1^{\beta\lambda} = 1$$

$$\langle 2 \rangle$$
7. Case:  $\alpha > 1$ 

PROOF:

$$(\alpha^{\beta})^{\lambda} = \sup_{\gamma < \lambda} (\alpha^{\beta})^{\gamma}$$

$$= \sup_{\gamma < \lambda} \alpha^{\beta\gamma}$$

$$= \alpha^{\sup_{\gamma < \lambda} \beta\gamma}$$

$$= \alpha^{\beta\lambda}$$

# Chapter 7

# Cardinal Numbers

## 7.1 Cardinal Numbers

**Definition 7.1.1** (Cardinality). For any set A, the *cardinality* or *cardinal number* |A| of A is the least ordinal equinumerous with A.

Let **Card** be the class of all cardinal numbers.

**Proposition 7.1.2.** For any sets A and B, we have  $A \approx B$  iff |A| = |B|.

Proof: Easy.  $\square$ 

**Definition 7.1.3** (Addition). Given cardinal numbers  $\kappa$  and  $\lambda$ , we define  $\kappa + \lambda$  to be  $|A \cup B|$  where A and B are disjoint sets of cardinality  $\kappa$  and  $\lambda$  respectively. We prove this is well-defined.

#### Proof:

- $\langle 1 \rangle 1$ . Assume:  $A \approx A'$ ,  $B \approx B'$ , and  $A \cap B = A' \cap B' = \emptyset$
- $\langle 1 \rangle 2$ . Pick bijections  $f: A \approx A'$  and  $g: B \approx B'$
- $\langle 1 \rangle 3$ . The function  $A \cup B \to A' \cup B'$  that maps  $a \in A$  to f(a) and  $b \in B$  to g(b) is a bijection.

**Proposition 7.1.4.** For any cardinal number  $\kappa$ , we have  $\kappa + 0 = \kappa$ .

PROOF: Let A and B be disjoint sets of cardinality  $\kappa$  and A. Then  $A = \emptyset$  so  $A \cup B = A$  and so  $A \cup B = \kappa$ .  $A \cap B = \emptyset$ 

**Theorem 7.1.5** (Associative Law for Addition). For any cardinal numbers  $\kappa$ ,  $\lambda$ ,  $\mu$  we have  $\kappa + (\lambda + \mu) = (\kappa + \lambda) + \mu$ .

PROOF: Since  $A \cup (B \cup C) = (A \cup B) \cup C$ .

**Proposition 7.1.6.** For any cardinal numbers  $\kappa$  and  $\lambda$  we have  $\kappa + \lambda = \lambda + \kappa$ .

PROOF: Since  $A \cup B = B \cup A$ .  $\square$ 

**Definition 7.1.7** (Multiplication). For  $\kappa$  and  $\lambda$  cardinal numbers, we define  $\kappa\lambda$  to be the cardinal number of  $A\times B$ , where  $|A|=\kappa$  and  $|B|=\lambda$ .

We prove this is well-defined.

PROOF: If  $f: A \approx A'$  and  $g: B \approx B'$  then the function that maps (a,b) to (f(a),g(b)) is a bijection  $A \times B \approx A' \times B'$ .  $\square$ 

**Proposition 7.1.8.** For any cardinal number  $\kappa$  we have  $\kappa \cdot 0 = 0$ .

PROOF: Since  $A \times \emptyset = \emptyset$ .  $\square$ 

**Proposition 7.1.9.** For any cardinal number  $\kappa$  we have  $\kappa \cdot 1 = \kappa$ .

PROOF: The function that maps (a, e) to a is a bijection  $A \times \{e\} \approx A$ .  $\square$ 

**Theorem 7.1.10** (Distributive Law). For any cardinal numbers  $\kappa$ ,  $\lambda$  and  $\mu$ , we have  $\kappa(\lambda + \mu) = \kappa\lambda + \kappa\mu$ .

PROOF: Since  $A \times (B \cup C) = (A \times B) \cup (A \times C)$ .  $\square$ 

**Theorem 7.1.11** (Associative Law for Multiplication). For any cardinal numbers  $\kappa$ ,  $\lambda$  and  $\mu$ , we have  $\kappa(\lambda\mu) = (\kappa\lambda)\mu$ .

PROOF: Since  $A \times (B \times C) \approx (A \times B) \times C$ .  $\square$ 

**Theorem 7.1.12** (Commutative Law for Multiplication). For any cardinal numbers  $\kappa$  and  $\lambda$ , we have  $\kappa\lambda = \lambda\kappa$ .

PROOF: Since  $A \times B \approx B \times A$ .  $\square$ 

**Theorem 7.1.13.** For any cardinal numbers  $\kappa$  and  $\lambda$ , if  $\kappa\lambda = 0$  then  $\kappa = 0$  or  $\lambda = 0$ .

PROOF: if  $A \times B = \emptyset$  then  $A = \emptyset$  or  $B = \emptyset$ .  $\square$ 

**Definition 7.1.14** (Exponentiation). Given cardinal numbers  $\kappa$  and  $\lambda$ , we define  $\kappa^{\lambda}$  to be  $|A^{B}|$ , where  $|A| = \kappa$  and  $|B| = \lambda$ .

We prove this is well-defined.

PROOF:If  $f: A \approx A'$  and  $g: B \approx B'$ , then the function that maps  $h: B \to A$  to  $f \circ h \circ g^{-1}$  is a bijection  $A^B \approx A'^{B'}$ .  $\square$ 

**Proposition 7.1.15.** For any cardinal numbers  $\kappa$ ,  $\lambda$  and  $\mu$ ,

$$\kappa^{\lambda+\mu} = (\kappa^{\lambda})^{\mu}$$

PROOF: The function that maps  $f: A \times B \to C$  to  $\lambda a \in A.\lambda b \in B.f(a,b)$  is a bijection  $A^{B \times C} \approx (A^B)^C$ .  $\square$ 

**Proposition 7.1.16.** For any cardinal numbers  $\kappa$ ,  $\lambda$  and  $\mu$ ,

$$(\kappa\lambda)^{\mu} = \kappa^{\mu}\lambda^{\mu}$$
.

PROOF: The function  $f: A^C \times B^C \to (A \times B)^C$  with f(g,h)(c) = (g(c),h(c)) is a bijection.  $\square$ 

**Proposition 7.1.17.** For any cardinal numbers  $\kappa$ ,  $\lambda$  and  $\mu$ , we have

$$\kappa^{\lambda+\mu} = \kappa^{\lambda} \kappa^{\mu} .$$

PROOF: If  $B \cap C = \emptyset$ , then  $f: A^B \times A^C \to A^{B \cup C}$  given by f(g,h)(b) = g(b) and f(g,h)(c) = h(c) is a bijection.  $\square$ 

**Proposition 7.1.18.** For any cardinal number  $\kappa$ , we have  $\kappa^0 = 1$ .

PROOF: For any set A, we have  $A^{\emptyset} = \{\emptyset\}$ .  $\square$ 

**Proposition 7.1.19.** For any cardinal number  $\kappa$ , we have  $\kappa^1 = \kappa$ .

PROOF: For any sets A and B, if  $B = \{b\}$  then the function  $f: A \to A^B$  with f(a)(b) = a is a bijection.  $\square$ 

**Proposition 7.1.20.** For any non-zero cardinal number  $\kappa$  we have  $0^{\kappa} = 0$ .

PROOF: If A is nonempty then there is no function  $A \to \emptyset$ .  $\square$ 

**Proposition 7.1.21.** For any set A we have  $|\mathcal{P}A| = 2^{|A|}$ .

PROOF: The function  $f: \mathcal{P}A \to 2^A$  where f(X)(a) = 0 if  $a \notin X$  and f(X)(a) = 1 if  $a \in X$ .  $\square$ 

Corollary 7.1.21.1. For any cardinal number  $\kappa$  we have  $\kappa < 2^{\kappa}$ .

PROOF: By Cantor's Theorem.

## 7.2 Ordering on Cardinal Numbers

**Definition 7.2.1.** Given cardinal numbers  $\kappa$  and  $\lambda$ , we have  $\kappa \leq \lambda$  iff  $A \leq B$ , where  $|A| = \kappa$  and  $|B| = \lambda$ .

### Proof:

- $\langle 1 \rangle 1$ . Let:  $|A| = \kappa$  and  $|B| = \lambda$
- $\langle 1 \rangle 2$ . Pick bijections  $f: A \approx \kappa$  and  $g: B \approx \lambda$
- $\langle 1 \rangle 3$ . If  $\kappa \leq \lambda$  then  $A \preccurlyeq B$

PROOF: Let  $i : \kappa \hookrightarrow \lambda$  be the inclusion. Then  $g^{-1} \circ i \circ f$  is an injection  $A \to B$ .

- $\langle 1 \rangle 4$ . If  $A \leq B$  then  $\kappa \leq \lambda$ 
  - $\langle 2 \rangle 1$ . Assume:  $A \leq B$
  - $\langle 2 \rangle 2$ . Pick an injection  $h: A \rightarrow B$
  - $\langle 2 \rangle 3.$   $g(h(A)) \subseteq B$  is well-ordered by  $\in$
  - $\langle 2 \rangle 4$ . Let:  $\gamma$  be the ordinal number of  $(g(h(A)), \in)$
  - $\langle 2 \rangle 5. \ \gamma \leq \lambda$

Proof: Proposition 6.1.12.

 $\langle 2 \rangle 6. \ \kappa \leq \gamma$ 

PROOF: By the leastness of  $\kappa$ , since A is equinumerous with  $\gamma$ .  $\langle 2 \rangle 7. \ \kappa \leq \lambda$ 

Corollary 7.2.1.1. There is no largest cardinal number.

**Proposition 7.2.2.** For any cardinal numbers  $\kappa$ ,  $\lambda$ ,  $\mu$ , if  $\kappa \leq \lambda$  then  $\kappa + \mu \leq \lambda + \mu$ .

PROOF: If  $f: A \to B$  is injective, and  $A \cap C = B \cap C = \emptyset$ , then the function  $A \cup C \to B \cup C$  that maps a to f(a) and maps c to c is an injection.  $\square$ 

**Proposition 7.2.3.** For any cardinal numbers  $\kappa$ ,  $\lambda$ ,  $\mu$ , if  $\kappa \leq \lambda$  then  $\kappa \mu \leq \lambda \mu$ .

PROOF: If  $f:A\to B$  is injective, then the function  $A\times C\to B\times C$  that maps (a,c) to (f(a),c) is injective.  $\square$ 

**Proposition 7.2.4.** For any cardinal numbers  $\kappa$ ,  $\lambda$ ,  $\mu$ , if  $\kappa \leq \lambda$  then  $\kappa^{\mu} \leq \lambda^{\mu}$ .

PROOF: Given an injection  $f:A\to B$ , the function that maps  $A^C\to B^C$  that maps g to  $f\circ g$  is an injection.  $\square$ 

**Proposition 7.2.5.** For any cardinal numbers  $\kappa$ ,  $\lambda$ ,  $\mu$ , if  $\kappa \leq \lambda$  and  $\mu$  and  $\kappa$  are not both 0, then  $\mu^{\kappa} \leq \mu^{\lambda}$ .

### Proof:

- $\langle 1 \rangle 1$ . Let: A, B and C be sets with A and C not both empty.
- $\langle 1 \rangle 2$ . Let:  $f: A \to B$  be an injection.

Prove:  $C^A \preceq C^B$ 

 $\langle 1 \rangle 3$ . Case:  $C = \emptyset$ 

PROOF: Then  $A \neq \emptyset$  so  $C^A = \emptyset \preccurlyeq C^B$ .

- $\langle 1 \rangle 4$ . Case:  $C \neq \emptyset$ 
  - $\langle 2 \rangle 1$ . Pick  $c \in C$
  - (2)2. Let:  $g: C^A \to C^B$  be the function g(h)(f(a)) = h(a), g(h)(b) = c if  $b \notin f(A)$

 $\langle 2 \rangle 3$ . g is an injection.

**Proposition 7.2.6.** Let A be a set such that  $\forall X \in A | X | \leq \kappa$ . Then

$$\left|\bigcup \mathcal{A}\right| \leq |\mathcal{A}|\kappa \ .$$

### Proof:

- $\langle 1 \rangle 1$ . For  $X \in \mathcal{A}$ , choose a surjection  $f_X : \kappa \to X$ .
- $\langle 1 \rangle 2$ . Define  $g: \mathcal{A} \times \kappa \to \bigcup \mathcal{A}$  by  $g(X, \alpha) = f_X(\alpha)$
- $\langle 1 \rangle 3$ . g is surjective.

**Lemma 7.2.7.** The union of a set of cardinal numbers is a cardinal number.

Proof:

 $\langle 1 \rangle 1.$  Let: A be a set of cardinal numbers.

Prove:  $\bigcup A$  is the smallest ordinal equinumerous with  $\bigcup A$ 

 $\langle 1 \rangle 2$ . Let:  $\alpha < \bigcup A$ 

Prove:  $\alpha \not\approx \bigcup A$  $\langle 1 \rangle 3$ . PICK  $\kappa \in A$  such that  $\alpha < \kappa$ 

 $\langle 1 \rangle 4$ .  $\alpha \prec \kappa$ 

 $\langle 1 \rangle 5. \ \alpha \prec \bigcup^{n} A$ 

# **Natural Numbers**

### 8.1 Inductive Sets

**Definition 8.1.1** (Inductive). A set I is *inductive* iff  $0 \in I$  and  $\forall x \in I.x^+ \in I$ .

**Definition 8.1.2** (Natural Number). A *natural number* is a set that belongs to every inductive set.

**Theorem 8.1.3.** The class  $\mathbb{N}$  of natural numbers is a set.

```
Proof: \langle 1 \rangle 1. Pick an inductive set I. Proof: Axiom of Infinity. \langle 1 \rangle 2. \mathbb{N} \subseteq I
```

**Theorem 8.1.4.**  $\mathbb{N}$  is inductive, and is a subset of every other inductive set.

```
PROOF: \langle 1 \rangle 1. \mathbb{N} is inductive, and is a subset of even \langle 2 \rangle 1. \mathbb{N} is inductive. \langle 2 \rangle 1. \mathbb{N} is inductive. \langle 2 \rangle 1. 0 \in \mathbb{N}

PROOF: Since 0 is a member of every inductive set. \langle 2 \rangle 2. \forall n \in \mathbb{N}. n^+ \in \mathbb{N}

\langle 3 \rangle 1. Let: n \in \mathbb{N}

\langle 3 \rangle 1. Let: n \in \mathbb{N}

\langle 3 \rangle 2. Let: n \in \mathbb{N}

PROVE: n^+ \in I

\langle 3 \rangle 3. n \in I

PROOF: \langle 3 \rangle 1, \langle 3 \rangle 2

\langle 3 \rangle 4. n^+ \in I

PROOF: \langle 3 \rangle 2, \langle 3 \rangle 3

\langle 1 \rangle 2. \mathbb{N} is a subset of every inductive set.

PROOF: Immediate from definitions.
```

**Corollary 8.1.4.1** (Induction Principle for  $\mathbb{N}$ ). Any inductive subset of  $\mathbb{N}$  coincides with  $\mathbb{N}$ .

**Theorem 8.1.5.** Every natural number except 0 is the successor of some natural number.

Proof: Trivially by induction.  $\square$ 

Proposition 8.1.6. Every natural number is an ordinal.

Proof: By induction.  $\square$ 

**Proposition 8.1.7.**  $\mathbb{N}$  is a transitive set.

#### Proof:

- $\langle 1 \rangle 1. \ 0 \subseteq \mathbb{N}$
- $\langle 1 \rangle 2. \ \forall n \in \mathbb{N}. n \subseteq \mathbb{N} \Rightarrow n^+ \subseteq \mathbb{N}$
- $\langle 1 \rangle 3. \ \forall n \in \mathbb{N}. n \subseteq \mathbb{N}$

PROOF: From  $\langle 1 \rangle 1$  and  $\langle 1 \rangle 2$  by induction.

Corollary 8.1.7.1. N is an ordinal.

**Definition 8.1.8.** We define  $\omega = \mathbb{N}$ .

**Proposition 8.1.9** (Dependent Choice). Let A be a nonempty set and R a relation on A such that  $\forall x \in A. \exists y \in A. (y,x) \in R$ . Then there exists a function  $f: \mathbb{N} \to A$  such that  $\forall n \in \mathbb{N}. (f(n+1),f(n)) \in R$ .

### Proof:

- $\langle 1 \rangle 1$ . PICK a choice function F for A.
- $\langle 1 \rangle 2$ . Pick  $a \in A$
- $\langle 1 \rangle 3$ . Define  $f : \mathbb{N} \to A$  by f(0) = a and  $f(n+1) = F(\{y \in A \mid (y, f(n)) \in R\})$ .

**Theorem Schema 8.1.10.** For any classes A and R, the following is a theorem:

Assume **R** is a relation on **A** and, for all  $a \in \mathbf{A}$ , the class  $\{x \in \mathbf{A} \mid x\mathbf{R}a\}$  is a set. Then **R** is well founded if and only if there does not exist a function  $f: \mathbb{N} \to \mathbf{A}$  such that  $\forall n \in \mathbb{N}. f(n+1)\mathbf{R}f(n)$ .

#### Proof:

 $\langle 1 \rangle 1$ . If there exists a function  $f : \mathbb{N} \to \mathbf{A}$  such that  $\forall n \in \mathbb{N}. f(n+1)\mathbf{R}f(n)$  then  $\mathbf{R}$  is not well founded.

PROOF:  $f(\mathbb{N})$  is a nonempty subset of **A** with no **R**-minimal element.

- $\langle 1 \rangle$ 2. If **R** is not well founded then there exists a function  $f : \mathbb{N} \to \mathbf{A}$  such that  $\forall n \in \mathbb{N}. f(n+1)\mathbf{R}f(n)$ .
  - $\langle 2 \rangle 1$ . Assume: **R** is not well founded.
  - $\langle 2 \rangle 2$ . Pick a nonempty subset  $B \subseteq \mathbf{A}$  that has no **R**-minimal element.
  - $\langle 2 \rangle 3. \ \forall x \in B. \exists y \in B. y \mathbf{R} x$

```
\langle 2 \rangle 4. Choose a function g: B \to B such that \forall x \in B.g(x)\mathbf{R}x \langle 2 \rangle 5. PICK b \in B \langle 2 \rangle 6. Define f: \mathbb{N} \to \mathbf{A} recursively by f(0) = b and \forall n \in \mathbb{N}.f(n+1) = g(f(n)) \langle 2 \rangle 7. \forall n \in \mathbb{N}.f(n+1)\mathbf{R}f(n)
```

## 8.2 Cardinality

**Definition 8.2.1** (Finite). A set is *finite* iff it is equinumerous to some natural number; otherwise it is *infinite*.

**Theorem 8.2.2** (Pigeonhole Principle). No natural number is equinumerous to a proper subset of itself.

```
PROOF: \langle 1 \rangle 1. Let: P(n) be the property: any one-to-one function n \to n is surjective. \langle 1 \rangle 2. P(0)
PROOF: The only function 0 \to 0 is injective.
```

 $\langle 1 \rangle 3$ . For every natural number n, if P(n) then P(n+1).

 $\langle 2 \rangle 1$ . Assume: P(n)

 $\langle 2 \rangle 2$ . Let: f be a one-to-one function  $n+1 \to n+1$ 

 $\langle 2 \rangle 3$ .  $f \upharpoonright n$  is a one-to-one function  $n \to n+1$ 

```
\langle 2 \rangle4. CASE: n \notin ranf

\langle 3 \rangle1. f \upharpoonright n : n \to n

\langle 3 \rangle2. ran(f \upharpoonright n) = n

\langle 3 \rangle3. f(n) = n

PROOF: \langle 2 \rangle1.

\langle 3 \rangle4. ran f = n + 1
```

 $\langle 2 \rangle 5$ . Case:  $n \in \operatorname{ran} f$ 

 $\langle 3 \rangle 1$ . Pick  $p \in n$  such that f(p) = n

 $(5/1. \text{ I lock } p \in n \text{ such that } f(p) = n$ 

 $\langle 3 \rangle 2$ . Let:  $\hat{f}: n \to n$  be the function

$$\hat{f}(p) = f(n)$$

$$\hat{f}(x) = f(x) \qquad (x \neq p)$$

 $\langle 3 \rangle 3$ .  $\hat{f}$  is one-to-one

 $\langle 3 \rangle 4$ . ran  $\hat{f} = n$ PROOF:  $\langle 2 \rangle 1$ 

 $\langle 3 \rangle$ 5. ran f = n + 1 $\langle 1 \rangle$ 4. For every natural number n, P(n).

Corollary 8.2.2.1. No finite set is equinumerous to a proper subset of itself.

Corollary 8.2.2.2. Every natural number is a cardinal number.

PROOF: For any natural number n, we have that n is the least ordinal such that  $n \approx n$ .  $\square$ 

Corollary 8.2.2.3.  $\mathbb{N}$  is a cardinal number.

Corollary 8.2.2.4.  $\mathbb{N}$  is infinite.

PROOF: The function that maps n to n+1 is a bijection between  $\mathbb N$  and  $\mathbb N-\{0\}$ .  $\square$ 

**Corollary 8.2.2.5.** If C is a proper subset of a natural number n, then there exists m < n such that  $C \approx m$ .

Proof: By Proposition 6.1.12.  $\square$ 

Corollary 8.2.2.6. Any subset of a finite set is finite.

**Proposition 8.2.3.** For any natural numbers m and n we have m+n (cardinal addition) is a natural number.

PROOF: Induction on n.  $\square$ 

Corollary 8.2.3.1. The union of two finite sets is finite.

Corollary 8.2.3.2. The union of a finite set of finite sets is finite.

Proof: By induction on the number of elements.  $\square$ 

**Proposition 8.2.4.** For natural numbers m and n, the cardinal sum m + n is equal to the ordinal sum m + n.

Proof: Induction on n.

**Proposition 8.2.5.** For any natural numbers m and n, we have mn (cardinal multiplication) is a natural number.

**Corollary 8.2.5.1.** If A and B are finite sets then  $A \times B$  is finite.

**Proposition 8.2.6.** For natural numbers m and n, the cardinal product mn is equal to the ordinal product mn.

Proof: Induction on n.

**Proposition 8.2.7.** For any natural numbers m and n we have  $m^n$  (cardinal exponentiation) is a natural number.

PROOF: Induction on n.

Corollary 8.2.7.1. If A and B are finite sets then  $A^B$  are finite.

**Proposition 8.2.8.** For natural numbers m and n, the cardinal exponentiation  $m^n$  and the ordinal exponentiation  $m^n$  agree.

PROOF: Induction on n.  $\square$ 

Proposition 8.2.9.  $\mathbb{N}^2 \approx \mathbb{N}$ 

PROOF: The function  $J: \mathbb{N}^2 \to \mathbb{N}$  defined by  $J(m,n) = ((m+n)^2 + 3m + n)/2$  is a bijection.  $\square$ 

**Proposition 8.2.10.** For any infinite cardinal  $\kappa$  we have  $\aleph_0 \leq \kappa$ .

#### Proof:

 $\langle 1 \rangle 1$ . Let: A be an infinite set.

Prove:  $\mathbb{N} \preceq A$ 

 $\langle 1 \rangle 2$ . PICK a choice function F for A.

 $\langle 1 \rangle 3$ . Define  $h : \mathbb{N} \to \{ X \in \mathcal{P}A \mid X \text{ is finite} \}$  by  $h(0) = \emptyset$ 

$$h(n+1) = h(n) \cup \{F(A - \{h(m) \mid m < n\})\}\$$

 $\langle 1 \rangle 4$ . Define  $g: \mathbb{N} \to A$  by  $g(n) = F(A - \{h(m) \mid m < n\})$ 

 $\langle 1 \rangle 5$ . g is injective.

PROOF: If m < n then  $g(m) \neq g(n)$ .

**Theorem Schema 8.2.11** (König's Lemma). For any classes **A** and **R**, the following is a theorem:

Assume **R** is a well founded relation on **A** such that, for all  $y \in \mathbf{A}$ , the class  $\{x \in \mathbf{A} \mid x\mathbf{R}y\}$  is a finite set. Let  $\mathbf{R}^t$  be the transitive closure of **R**. Then, for all  $y \in \mathbf{A}$ , the class  $\{x \in \mathbf{A} \mid x\mathbf{R}^ty\}$  is a finite set.

### Proof:

 $\langle 1 \rangle 1$ . Let:  $y \in \mathbf{A}$ 

 $\langle 1 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall x \mathbf{R} y . \{z \in \mathbf{A} \mid z \mathbf{R}^t x\}$  is a finite set.

 $\langle 1 \rangle 3. \ \{x \mid x\mathbf{R}^ty\} = \bigcup_{x\mathbf{R}y} (\{x\} \cup \{z \mid z\mathbf{R}^tx\}$ 

 $\langle 1 \rangle 4$ .  $\{x \mid x \mathbf{R}^t y\}$  is finite.

Proof: Corollary 8.2.3.2.

## 8.3 Countable Sets

**Definition 8.3.1** (Countable). A set A is countable iff  $|A| \leq \aleph_0$ .

**Theorem 8.3.2.** The union of a countable set of countable sets is countable.

Proof: Proposition 7.2.6.  $\square$ 

### 8.4 Arithmetic

**Definition 8.4.1** (Even). A natural number n is *even* iff there exists  $m \in \mathbb{N}$  such that n = 2m.

**Definition 8.4.2** (Odd). A natural number n is odd iff there exists  $p \in \mathbb{N}$  such that n = 2p + 1.

```
Proposition 8.4.3. Every natural number is either even or odd.
Proof:
\langle 1 \rangle 1. 0 is even.
   Proof: 0 = 2 \times 0.
\langle 1 \rangle 2. For any natural number n, if n is either even or odd then n^+ is either even
   Proof:
   \langle 2 \rangle 1. Let: n \in \mathbb{N}
   \langle 2 \rangle 2. If n is even then n^+ is odd.
     PROOF: If n = 2p then n^+ = 2p + 1.
   \langle 2 \rangle 3. If n is odd then n^+ is even.
      PROOF: If n = 2p + 1 then n^{+} = 2(p + 1).
Proposition 8.4.4. No natural number is both even and odd.
Proof:
\langle 1 \rangle 1. 0 is not odd.
   PROOF: For any p we have 2p + 1 = (2p)^+ \neq 0.
\langle 1 \rangle 2. For any natural number n, if n is not both even and odd, then n^+ is not
       both even and odd.
   \langle 2 \rangle 1. Let: n be a natural number.
   \langle 2 \rangle 2. If n^+ is even then n is odd.
      \langle 3 \rangle 1. Assume: n^+ is even.
      \langle 3 \rangle 2. PICK p such that n^+ = 2p
      \langle 3 \rangle 3. \ p \neq 0
        PROOF: Since n^+ \neq 0.
```

 $\langle 3 \rangle 4$ . Pick q such that  $p = q^+$ 

PROOF: Theorem 8.1.5.

 $\langle 3 \rangle 5. \ n^+ = 2q + 2$ 

Proof:  $\langle 3 \rangle 2$ ,  $\langle 3 \rangle 4$ .

 $\langle 3 \rangle 6. \ \ n = 2q + 1$ 

Proof: Proposition 6.2.7,  $\langle 3 \rangle 5$ 

 $\langle 3 \rangle 7$ . *n* is odd.

 $\langle 2 \rangle 3$ . If  $n^+$  is odd then n is even.

 $\langle 3 \rangle 1$ . Assume:  $n^+$  is odd.

 $\langle 3 \rangle 2$ . Pick p such that  $n^+ = 2p + 1$ 

 $\langle 3 \rangle 3$ . n = 2p

Proof: Proposition 6.2.7,  $\langle 3 \rangle 2$ 

 $\langle 3 \rangle 4$ . *n* is even.

**Proposition 8.4.5.** Let m, n, p, q be natural numbers. Assume m+n=p+q. Then m < p if and only if q < n.

#### Proof:

 $\langle 1 \rangle 1$ . If m < p then q < n.

PROOF: If m < p and  $n \le q$  then  $m + n . <math>\langle 1 \rangle 2$ . If q < n then m < p. PROOF: Similar.

**Proposition 8.4.6.** Let m, n, p and q be natural numbers. Assume n < m and q < p. Then

$$mq + np < mp + nq$$
.

Proof:

 $\langle 1 \rangle 1$ . Pick positive natural numbers a and b such that m=n+a and p=q+b.

 $\langle 1 \rangle 2$ . mp + nq > mq + np

Proof:

$$mp + nq = (n+a)(q+b) + nq$$

$$= 2nq + nb + aq + ab$$

$$mq + np = (n+a)q + n(q+b)$$

$$= 2nq + aq + nb$$

$$\therefore mp + nq = mq + np + ab$$

$$> mq + np$$

## 8.5 Sequences

**Definition 8.5.1** (Sequence). Let A be a set. A *finite sequence* in A is a function  $a:n\to A$  for some natural number n; we write it as  $(a(0),a(1),\ldots,a(n-1))$ . An *(infinite) sequence* in A is a function  $\mathbb{N}\to A$ .

We write  $A^*$  for the set of all finite sequences in A.

**Proposition 8.5.2.** If A is countable then  $A^*$  is countable.

PROOF: For any n, the set  $A^n$  is countable, and  $A^*$  is equinumerous with  $\bigcup_n A^n$ .

### 8.6 Transitive Closure of a Set

**Proposition 8.6.1.** For any set A, there exists a unique transitive set C such that:

- $A \subseteq C$
- For any transitive set X, if  $A \subseteq X$  then  $C \subseteq X$

Proof:

 $\langle 1\rangle 1.$  Define a function  $F:\mathbb{N}\to \mathbf{V}$  by F(0)=A  $F(n+1)=A\cup \bigcup (F(0)\cup\cdots\cup F(n))$ 

```
\langle 1 \rangle 2. For all n \in \mathbb{N} and a \in F(n) we have a \subseteq F(n+1)
    PROOF: a \in F(0) \cup \cdots \cup F(n) so a \subseteq \bigcup (F(0) \cup \cdots \cup F(n)) \subseteq F(n+1).
\langle 1 \rangle 3. Let: C = \bigcup_{n \in \mathbb{N}} F(n)
\langle 1 \rangle 4. C is transitive.
    \langle 2 \rangle 1. Let: x \in y \in C
    \langle 2 \rangle 2. Pick n \in \mathbb{N} such that y \in F(n)
    \langle 2 \rangle 3. \ y \subseteq F(n+1)
       Proof: \langle 1 \rangle 2
    \langle 2 \rangle 4. \ x \in F(n+1)
    \langle 2 \rangle 5. \ x \in C
\langle 1 \rangle 5. A \subseteq C
    PROOF: Since F(0) = A.
\langle 1 \rangle 6. For any transitive set X, if A \subseteq X then C \subseteq X
    \langle 2 \rangle 1. Let: X be a transitive set
    \langle 2 \rangle 2. Assume: A \subseteq X
    \langle 2 \rangle 3. For all n \in \mathbb{N} we have F(n) \subseteq X.
        \langle 3 \rangle 1. \ F(0) \subseteq X
            Proof: \langle 2 \rangle 2
        \langle 3 \rangle 2. For all n \in \mathbb{N}, if F(n) \subseteq X, then F(n+1) \subseteq X.
            \langle 4 \rangle 1. Let: n \in \mathbb{N}
            \langle 4 \rangle 2. Assume: \forall m < n.F(m) \subseteq X
            \langle 4 \rangle 3. \ F(0) \cup \cdots \cup F(n) \subseteq X
            \langle 4 \rangle 4. \bigcup (F(0) \cup \cdots \cup F(n)) \subseteq X
               Proof: Since X is transitive.
            \langle 4 \rangle 5. F(n+1) \subseteq X
    \langle 2 \rangle 4. \ C \subseteq X
\langle 1 \rangle 7. Let D be a transitive set such that A \subseteq D and, for any transitive set X,
          if A \subseteq X then D \subseteq X. Then D = C.
    PROOF: We have C \subseteq D and D \subseteq C.
```

### 8.7 The Veblen Fixed Point Theorem

**Theorem Schema 8.7.1** (Veblen Fixed Point Theorem). For any class **T**, the following is a theorem:

Assume **T** is a normal ordinal operation. For every ordinal  $\beta$ , there exists  $\gamma \geq \beta$  such that  $\mathbf{T}(\gamma) = \gamma$ .

### Proof:

 $\gamma := \beta$ .

- $\langle 1 \rangle 1$ . Let:  $\beta$  be an ordinal.
- $\langle 1 \rangle$ 2. Assume: w.l.o.g.  $\beta < \mathbf{T}(\beta)$ PROOF: We have  $\beta \leq \mathbf{T}(\beta)$  by Proposition 6.4.5, and if  $\beta = \mathbf{T}(\beta)$  we take

 $\langle 1 \rangle 3$ . Define  $f : \mathbb{N} \to \mathbf{On}$  by recursion thus:

$$f(0) = \beta$$

$$f(n^{+}) = \mathbf{T}(f(n))$$

$$\langle 1 \rangle 4. \text{ Let: } \gamma = \sup_{n \in \mathbb{N}} f(n)$$

$$\langle 1 \rangle 5. \beta \leq \gamma$$

$$\text{Proof: Since } \beta = f(0).$$

$$\langle 1 \rangle 6. \mathbf{T}(\gamma) = \gamma$$

$$\langle 2 \rangle 1. \mathbf{T}(\gamma) \leq \gamma$$

$$\text{Proof:}$$

$$\mathbf{T}(\gamma) = \sup_{n \in \mathbb{N}} \mathbf{T}(f(n)) \qquad \text{(Theorem 6.4.7)}$$

$$= \sup_{n \in \mathbb{N}} f(n^{+}) \qquad \text{($\langle 1 \rangle 3$)}$$

$$\leq \sup_{n \in \mathbb{N}} f(n)$$

$$= \gamma$$

$$\langle 2 \rangle 2. \gamma \leq \mathbf{T}(\gamma)$$

$$\text{Proof: Proposition 6.4.5.}$$

**Definition 8.7.2** (Derived Operation). Let T be a normal ordinal operation. The *derived* operation  $T': On \to V$  is the unique order isomorphism between On and the fixed points of T.

**Proposition Schema 8.7.3.** For any class  $\mathbf{T}$ , the following is a theorem: If  $\mathbf{T}$  is a normal ordinal operation, then the derived operation is normal.

Proof:

- $\langle 1 \rangle 1$ . For any set S of fixed points of **T**, we have  $\bigcup S$  is a fixed point of **T**  $\langle 2 \rangle 1$ . LET: S be a set of fixed points of **T**.
  - $\langle 2 \rangle 2$ .  $\mathbf{T}(\sup S) = \sup S$

Proof:

$$\begin{aligned} \mathbf{T}(\sup S) &= \sup_{\alpha \in S} \mathbf{T}(\alpha) & \text{(Theorem 6.4.7)} \\ &= \sup_{\alpha \in S} \alpha & \text{($\langle 2 \rangle 1$)} \\ &= \sup S \end{aligned}$$

 $\langle 1 \rangle 2$ . Q.E.D.

Proof: Proposition 6.4.8.

### 8.8 Cantor Normal Form

**Theorem 8.8.1.** For any ordinal  $\alpha$ , there exist a unique sequence of nonzero natural numbers  $(n_1, \ldots, n_k)$  and sequence of ordinals  $(\gamma_1, \ldots, \gamma_k)$  such that

$$\gamma_k < \gamma_{k-1} < \dots < \gamma_1$$

and

$$\alpha = \omega^{\gamma_1} n_1 + \omega^{\gamma_2} n_2 + \dots + \omega^{\gamma_k} n_k .$$

Proof:

 $\langle 1 \rangle 1$ . For any ordinal  $\alpha$ , there exist a sequence of nonzero natural numbers  $(n_1, \ldots, n_k)$  and sequence of ordinals  $(\gamma_1, \ldots, \gamma_k)$  such that

$$\gamma_k < \gamma_{k-1} < \dots < \gamma_1$$

and

$$\alpha = \omega^{\gamma_1} n_1 + \omega^{\gamma_2} n_2 + \dots + \omega^{\gamma_k} n_k .$$

- $\langle 2 \rangle 1$ . Let:  $\alpha$  be an ordinal
- $\langle 2 \rangle 2$ . Assume: as an induction hypothesis that, for all  $\beta < \alpha$ , the theorem holds.
- $\langle 2 \rangle 3$ . Assume: w.l.o.g.  $\alpha \neq 0$
- $\langle 2 \rangle 4$ . Let:  $\gamma_1$ ,  $n_1$ ,  $\rho_1$  be the unique ordinals such that  $0 \neq n_1 < \omega$ ,  $\rho_1 < \omega^{\gamma_1}$ , and  $\alpha = \omega^{\gamma_1} n_1 + \rho_1$
- $\langle 2 \rangle$ 5. Let:  $(\gamma_2, \dots, \gamma_k)$  and  $(n_2, \dots, n_k)$  be sequences such that  $\gamma_k < \gamma_{k-1} < \dots < \gamma_2$  and  $\rho_1 = \omega^{\gamma_2} n_2 + \dots + \omega^{\gamma_k} n_k$
- $\langle 2 \rangle 6. \ \gamma_2 < \gamma_1$

PROOF: Since  $\omega^{\gamma_2} \leq \rho_1 < \omega^{\gamma_1}$ 

 $\langle 1 \rangle 2$ . If

$$\gamma_k < \gamma_{k-1} < \dots < \gamma_1 \gamma'_k < \gamma'_{k-1} < \dots < \gamma'_1$$

and

$$\omega^{\gamma_1} n_1 + \omega^{\gamma_2} n_2 + \dots + \omega^{\gamma_k} n_k = \omega^{\gamma'_1} n'_1 + \omega^{\gamma'_2} n'_2 + \dots + \omega^{\gamma'_k} n'_k$$
then  $\gamma_i = \gamma'_i$  for all  $i$  and  $n_i = n'_i$  for all  $i$ 

PROOF: Prove by induction on i using the Logarithm Theorem.

**Definition 8.8.2** (Cantor Normal Form). For any ordinal  $\alpha$ , the *Cantor normal* form of  $\alpha$  is the expression  $\alpha = \omega^{\gamma_1} n_1 + \cdots + \omega^{\gamma_k} n_k$  such that  $n_1, \ldots, n_k$  are nonzero natural numbers and  $\gamma_k < \gamma_{k-1} < \cdots < \gamma_1$ .

# **Infinite Cardinals**

## 9.1 Alephs

**Definition 9.1.1** (Aleph). Let  $\aleph$  be the unique order isomorphism between **On** and the class of infinite cardinals.

**Proposition 9.1.2.** The operation  $\aleph$  is normal.

Proof: Proposition 6.4.8 and Lemma 7.2.7.  $\Box$ 

**Definition 9.1.3** (Continuum Hypothesis). The *continuum hypothesis* is the statement that  $\aleph_1 = 2^{\aleph_0}$ .

**Definition 9.1.4** (Generalised Continuum Hypothesis). The *generalised continuum hypothesis* is the statement that, for all  $\alpha$ ,  $\aleph_{\alpha^+} = 2^{\aleph_{\alpha}}$ .

### 9.2 Beths

**Definition 9.2.1** (Beth). Define the operation  $\beth: \mathbf{On} \to \mathbf{Card}$  by transfinite recursion as follows:

$$\begin{split} & \beth_0 := \aleph_0 \\ & \beth_{\alpha^+} := 2^{\beth_\alpha} \\ & \beth_\lambda := \bigcup_{\alpha < \lambda} \beth_\alpha \end{split} \qquad \qquad (\lambda \text{ a limit ordinal})$$

**Proposition 9.2.2.**  $\supset$  *is a normal operation.* 

PROOF: It is continuous by definition, and  $\beth_{\alpha} < \beth_{\alpha^+}$  by Cantor's Theorem.  $\square$ 

**Proposition 9.2.3.** The continuum hypothesis is equivalent to the statement  $\beth_1 = \aleph_1$ .

The generalised continuum hypothesis is equivalent to the statement  $\beth = \alpha$ .

PROOF: Immediate from definitions.

# The Cumulative Hierarchy

**Definition 10.0.1** (Cumulative Hierarchy). Define the function  $V: \mathbf{On} \to \mathbf{V}$  by transfinite recursion thus:

$$V_{\alpha} = \bigcup_{\beta < \alpha} \mathcal{P}V_{\beta}$$

**Proposition 10.0.2.** For all  $\alpha \in \mathbf{On}$ ,  $V_{\alpha}$  is a transitive set.

Proof:

 $\langle 1 \rangle 1$ . Let:  $\alpha \in \mathbf{On}$ 

 $\langle 1 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall \beta < \alpha. V_{\beta}$  is a transitive set.

 $\langle 1 \rangle 3$ . For all  $\beta < \alpha$ ,  $\mathcal{P}V_{\beta}$  is a transitive set.

PROOF: Proposition 1.6.4.  $\langle 1 \rangle 4$ .  $V_{\alpha}$  is a transitive set. PROOF: Proposition 1.6.3.

ш

**Proposition 10.0.3.** For any ordinals  $\alpha$  and  $\beta$ , if  $\beta < \alpha$  then  $V_{\beta} \subseteq V_{\alpha}$ .

PROOF: Since  $V_{\beta} \in \mathcal{P}V_{\beta} \subseteq V_{\alpha}$  and  $V_{\alpha}$  is a transitive set.  $\square$ 

Theorem 10.0.4.

1. 
$$V_0 = \emptyset$$

2. 
$$\forall \alpha \in \mathbf{On}.V_{\alpha^+} = \mathcal{P}V_{\alpha}$$

3. For any limit ordinal  $\lambda$ ,  $V_{\lambda} = \bigcup_{\alpha \leq \lambda} V_{\alpha}$ .

Proof:

 $\langle 1 \rangle 1. \ V_0 = \emptyset$ 

PROOF: Immediate from definition.

 $\langle 1 \rangle 2. \ \forall \alpha \in \mathbf{On}.V_{\alpha^+} = \mathcal{P}V_{\alpha}$ 

Proof:

- $\langle 2 \rangle 1$ . Let:  $\alpha \in \mathbf{On}$
- $\langle 2 \rangle 2$ . For all  $\beta < \alpha$  we have  $\mathcal{P}V_{\beta} \subseteq \mathcal{P}V_{\alpha}$ PROOF: Propositions 1.5.8 and 10.0.3.
- $\langle 2 \rangle 3. \ V_{\alpha^+} = \mathcal{P} V_{\alpha}$

$$V_{\alpha^{+}} = \bigcup_{\beta < \alpha^{+}} \mathcal{P}V_{\beta}$$

$$= \bigcup_{\beta < \alpha} \mathcal{P}V_{\beta} \cup \mathcal{P}V_{\alpha}$$

$$\mathcal{P}V$$

 $\langle 1 \rangle$ 3. For any limit ordinal  $\lambda$ ,  $V_{\lambda} = \bigcup_{\alpha < \lambda} V_{\alpha}$ 

Proof:

 $\langle 2 \rangle 1. \ V_{\lambda} \subseteq \bigcup_{\alpha < \lambda} V_{\alpha}$ 

Proof:

$$V_{\lambda} = \bigcup_{\alpha < \lambda} \mathcal{P}V_{\alpha}$$

$$= \bigcup_{\alpha < \lambda} V_{\alpha^{+}} \qquad (\langle 1 \rangle 2)$$

$$\subseteq \bigcup_{\alpha < \lambda} V_{\alpha}$$

 $\langle 2 \rangle 2. \bigcup_{\alpha < \lambda} V_{\alpha} \subseteq V_{\lambda}$ PROOF: Proposition 10.0.3.

**Proposition 10.0.5.** For every set A, there exists an ordinal  $\alpha$  such that  $A \in$  $V_{\alpha}$ .

### Proof:

- $\langle 1 \rangle 1$ . Let us say a set A is grounded iff there exists an ordinal  $\alpha$  such that  $A \in V_{\alpha}$ .
- $\langle 1 \rangle 2$ . For any set A, if every element of A is grounded, then A is grounded.
  - $\langle 2 \rangle 1$ . Let: A be a set.
  - $\langle 2 \rangle 2$ .  $S = \{ \alpha \mid \exists a \in A.\alpha \text{ is the least ordinal such that } a \in V_{\alpha} \}$

PROOF: S is a set by an Axiom of Replacement.

- $\langle 2 \rangle 3$ . Let:  $\beta = \sup S$
- $\langle 2 \rangle 4$ .  $A \subseteq V_{\beta}$ 
  - $\langle 3 \rangle 1$ . Let:  $a \in A$
  - $\langle 3 \rangle 2$ . Let:  $\alpha$  be the least ordinal such that  $a \in V_{\beta}$
  - $\langle 3 \rangle 3. \ \alpha \in S$
  - $\langle 3 \rangle 4. \ \alpha \leq \beta$
  - $\langle 3 \rangle 5. \ a \in V_{\beta}$
- $\langle 2 \rangle 5. \ A \in V_{\beta^+}$
- $\langle 1 \rangle 3$ . Assume: for a contradiction there exists an ungrounded set.
- $\langle 1 \rangle 4$ . PICK a transitive set B that has an ungrounded member.

PROOF: Pick a transitive set c, and take B to be the transitive closure of  $\{c\}$ .

```
\begin{array}{l} \langle 1 \rangle 5. \text{ Let: } A = \{x \in B \mid x \text{ is ungrounded}\} \\ \langle 1 \rangle 6. \text{ Pick } m \in A \text{ such that } m \cap A = \emptyset \\ \text{Proof: Axiom of Regularity.} \\ \langle 1 \rangle 7. \text{ Every member of } m \text{ is grounded.} \\ \langle 2 \rangle 1. \text{ Assume: for a contradiction } x \in m \text{ is ungrounded.} \\ \langle 2 \rangle 2. x \in B \\ \text{Proof: Since } B \text{ is transitive } (\langle 1 \rangle 4). \\ \langle 2 \rangle 3. x \in A \\ \text{Proof: } \langle 1 \rangle 5 \\ \langle 2 \rangle 4. \text{ Q.E.D.} \\ \text{Proof: This contradicts } \langle 1 \rangle 6. \\ \langle 1 \rangle 8. m \text{ is grounded.} \\ \text{Proof: } \langle 1 \rangle 2 \\ \langle 1 \rangle 9. \text{ Q.E.D.} \\ \text{Proof: This contradicts } \langle 1 \rangle 6. \\ \end{array}
```

**Definition 10.0.6** (Rank). The *rank* of a set A is the least ordinal  $\alpha$  such that  $A \in V_{\alpha^+}$ .

**Proposition 10.0.7.** For any set A we have

$$\operatorname{rank} A = \bigcup_{a \in A} (\operatorname{rank} a)^+$$

```
PROOF:  \langle 1 \rangle 1. \text{ Let: } \alpha = \bigcup_{a \in A} (\operatorname{rank} a)^+   \langle 1 \rangle 2. \quad A \subseteq V_{\alpha}   \langle 2 \rangle 1. \text{ Let: } a \in A   \langle 2 \rangle 2. \quad a \in V_{(\operatorname{rank} a)^+}   \langle 2 \rangle 3. \quad a \in V_{\alpha}   \langle 1 \rangle 3. \quad A \in V_{\alpha^+}   \langle 1 \rangle 4. \text{ If } A \subseteq V_{\beta} \text{ then } \alpha \leq \beta   \langle 2 \rangle 1. \text{ Assume: } A \subseteq V_{\beta}   \langle 2 \rangle 2. \text{ For all } a \in A \text{ we have } (\operatorname{rank} a)^+ \leq \beta   \text{PROOF: Since } a \in V_{\beta}.   \langle 2 \rangle 3. \quad \alpha \leq \beta   \square
```

Corollary 10.0.7.1. For any sets a and b, if  $a \in b$  then rank  $a < \operatorname{rank} b$ .

**Proposition 10.0.8.** For any ordinal number  $\alpha$  we have rank  $\alpha = \alpha$ .

### Proof:

- $\langle 1 \rangle 1$ . Let:  $\alpha$  be an ordinal.
- $\langle 1 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall \beta < \alpha$ . rank  $\beta = \beta$
- $\langle 1 \rangle 3$ . rank  $\alpha = \bigcup_{\beta < \alpha} \beta^+$

Proof:

$$\operatorname{rank} \alpha = \bigcup_{\beta < \alpha} (\operatorname{rank} \beta)^+$$
$$= \bigcup_{\beta < \alpha} \beta^+$$

 $\begin{array}{c} \beta < \alpha \\ \langle 1 \rangle 4. \ \bigcup_{\beta < \alpha} \beta^+ \leq \alpha \\ \text{Proof: Since for all } \beta < \alpha \text{ we have } \beta^+ \leq \alpha. \\ \langle 1 \rangle 5. \ \alpha \leq \bigcup_{\beta < \alpha} \beta^+ \\ \langle 2 \rangle 1. \ \text{Let: } \gamma = \bigcup_{\beta < \alpha} \beta^+ \\ \langle 2 \rangle 2. \ \text{Assume: for a contradiction } \gamma < \alpha \\ \langle 2 \rangle 3. \ \gamma^+ \leq \bigcup_{\beta < \alpha} \beta^+ = \gamma \\ \langle 2 \rangle 4. \ \text{Q.E.D.} \\ \text{Proof: This is a contradiction} \end{array}$ 

PROOF: This is a contradiction.

# Models of Set Theory

**Definition 11.0.1** (Relativization). Let  $\sigma$  be a sentence in the language of set theory and  $\mathbf{M}$  a class. The *relativization* of  $\sigma$  to  $\mathbf{M}$  is the sentence  $\sigma^{\mathbf{M}}$  formed by replacing every quantifier  $\forall x$  with  $\forall x \in \mathbf{M}$ , and  $\exists x$  with  $\exists x \in \mathbf{M}$ .

We write 'M is a model of  $\sigma$ ' for the sentence  $\sigma^{\mathbf{M}}$ .

**Theorem Schema 11.0.2.** For any class M, the following is a theorem: If M is a transitive class, then M is a model of the Axiom of Extensionality.

```
Proof:
```

```
\begin{array}{ll} \langle 1 \rangle 1. \text{ Assume: } \mathbf{M} \text{ is a transitive class.} \\ \text{Prove: } \forall x,y \in \mathbf{M} (\forall z \in \mathbf{M} (z \in x \Leftrightarrow z \in y) \Rightarrow x = y) \\ \langle 1 \rangle 2. \text{ Let: } x,y \in \mathbf{M} \\ \langle 1 \rangle 3. \text{ Assume: } \forall z \in \mathbf{M} (z \in x \Leftrightarrow z \in y) \\ \langle 1 \rangle 4. \ \forall z (z \in x \Leftrightarrow z \in y) \\ \text{Proof: Since } z \in x \Rightarrow z \in \mathbf{M} \text{ and } z \in y \Rightarrow z \in \mathbf{M} \text{ by } \langle 1 \rangle 1. \\ \langle 1 \rangle 5. \ x = y \\ \square \end{array}
```

**Theorem 11.0.3.** If  $\alpha$  is a non-zero ordinal then  $V_{\alpha}$  is a model of the statement: The empty class is a set.

### Proof:

```
\begin{array}{ll} \langle 1 \rangle 1. & \text{Let: } \alpha \neq 0 \\ & \text{Prove: } \exists x \in V_{\alpha}. \forall y \in V_{\alpha}. y \notin x \\ \langle 1 \rangle 2. & \emptyset \in V_{\alpha} \\ \langle 1 \rangle 3. & \forall y \in V_{\alpha}. y \notin \emptyset \\ & \Box \end{array}
```

**Theorem 11.0.4.** For any limit ordinal  $\lambda$ , we have  $V_{\lambda}$  is a model of the statement: for any sets a and b, the class  $\{a,b\}$  is a set.

### Proof:

 $\langle 1 \rangle 1$ . Let:  $\lambda$  be a limit ordinal.

```
PROVE: \forall a,b \in V_{\lambda}. \exists c \in V_{\lambda}. \forall x \in V_{\lambda} (x \in c \Leftrightarrow x = a \lor x = b) \langle 1 \rangle 2. Let: a,b \in V_{\lambda} \langle 1 \rangle 3. Pick \alpha,\beta < \lambda such that a \in V_{\alpha} and b \in V_{\beta} \langle 1 \rangle 4. Assume: w.l.o.g. \alpha \leq \beta \langle 1 \rangle 5. a,b \in V_{\beta} \langle 1 \rangle 6. \{a,b\} \in V_{\beta+1} \langle 1 \rangle 7. \{a,b\} \in V_{\lambda} \langle 1 \rangle 8. \forall x \in V_{\lambda} (x \in \{a,b\} \Leftrightarrow x = a \lor x = b)
```

**Theorem 11.0.5.** For any ordinal  $\alpha$ , we have  $V_{\alpha}$  is a model of the Union Axiom.

```
Proof:
```

```
\begin{array}{l} \text{TROOF:} \\ \langle 1 \rangle 1. \text{ Let: } \alpha \text{ be an ordinal.} \\ \qquad \qquad \text{PROVE:} \quad \forall a \in V_{\alpha}. \exists b \in V_{\alpha}. \forall x \in V_{\alpha} (x \in b \Leftrightarrow \exists y \in V_{\alpha} (x \in y \land y \in a)) \\ \langle 1 \rangle 2. \text{ Let: } a \in V_{\alpha} \\ \langle 1 \rangle 3. \text{ PICK } \beta < \alpha \text{ such that } a \subseteq V_{\beta} \\ \langle 1 \rangle 4. \quad \bigcup a \subseteq V_{\beta} \\ \qquad \qquad \text{PROOF: } V_{\beta} \text{ is a transitive set.} \\ \langle 1 \rangle 5. \quad \bigcup a \in V_{\alpha} \\ \langle 1 \rangle 6. \quad \forall x \in V_{\alpha} (x \in \bigcup a \Leftrightarrow \exists y \in V_{\alpha} (x \in y \land y \in a)) \\ \qquad \qquad \qquad \text{PROOF: } V_{\alpha} \text{ is a transitive set.} \\ \Box \end{array}
```

**Theorem 11.0.6.** For any limit ordinal  $\lambda$ , we have  $V_{\lambda}$  is a model of the Power Set Axiom.

### Proof:

```
\begin{array}{l} \text{TROOT.} \\ \langle 1 \rangle \text{1. Let: } \lambda \text{ be a limit ordinal.} \\ \text{PROVE: } \forall a \in V_{\lambda}. \exists b \in V_{\lambda}. \forall x \in V_{\lambda} (x \in b \Leftrightarrow \forall y \in V_{\lambda} (y \in x \Rightarrow y \in a)) \\ \langle 1 \rangle \text{2. Let: } a \in V_{\lambda} \\ \langle 1 \rangle \text{3. PICK } \alpha < \lambda \text{ such that } a \in V_{\alpha} \\ \langle 1 \rangle \text{4. } \mathcal{P}a \in V_{\alpha+1} \\ \langle 1 \rangle \text{5. } \mathcal{P}a \in V_{\lambda} \\ \langle 1 \rangle \text{6. } \forall x \in V_{\lambda} (x \in \mathcal{P}a \Leftrightarrow \forall y \in V_{\lambda} (y \in x \Rightarrow y \in a)) \\ & \square \end{array}
```

**Theorem Schema 11.0.7.** For any property  $P[x, y_1, ..., y_n]$ , the following is a theorem:

For any ordinal  $\alpha$ , the set  $V_{\alpha}$  is a model of the statement: for any sets  $a_1$ , ...,  $a_n$ , B, the class  $\{x \in B \mid P[x, a_1, ..., a_n]\}$  is a set.

### Proof:

 $\langle 1 \rangle 1$ . Let:  $\alpha$  be an ordinal.  $\langle 1 \rangle 2$ . Let:  $a_1, \ldots, a_n, B \in V_{\alpha}$   $\langle 1 \rangle 3$ . Let:  $C = \{x \in B \mid P[x, a_1, \ldots, a_n]^{V_{\alpha}}\}$   $\langle 1 \rangle 4$ .  $C \in V_{\alpha}$ 

```
\langle 1 \rangle 5. \ \forall x \in V_{\alpha}(x \in C \Leftrightarrow x \in B \land P[x, a_1, \dots, a_n]^{V_{\alpha}})
```

**Theorem 11.0.8.** For any ordinal  $\alpha > \omega$ , we have:  $V_{\alpha}$  is a model of the Axiom of Infinity.

### Proof:

- $\langle 1 \rangle 1$ . Let:  $\alpha > \omega$
- $\langle 1 \rangle 2$ .  $\mathbb{N} \in V_{\alpha}$
- $\langle 1 \rangle 3. \ \exists e \in V_{\alpha}(e \in \mathbb{N} \land \forall x \in V_{\alpha}.x \notin e)$
- $\langle 1 \rangle 4. \ \forall x \in V_{\alpha}(x \in \mathbb{N} \Rightarrow \exists y \in V_{\alpha} \forall z \in V_{\alpha}(z \in y \Leftrightarrow z \in x \lor z = x))$

**Theorem 11.0.9.** For any ordinal  $\alpha$ , we have  $V_{\alpha}$  is a model of the Axiom of Choice.

### Proof:

- $\langle 1 \rangle 1$ . Let:  $\alpha$  be an ordinal.
- $\langle 1 \rangle 2$ . Let:  $A \in V_{\alpha}$
- $\langle 1 \rangle 3$ . Assume:  $\forall x \in V_{\alpha} (x \in A \Rightarrow \exists y \in V_{\alpha}.y \in A)$
- $\langle 1 \rangle 4$ . Assume:  $\forall x, y, z \in V_{\alpha} (x \in A \land y \in A \land z \in x \land z \in y \Rightarrow x = y)$
- $\langle 1 \rangle$ 5. A is a set of pairwise disjoint nonempty sets.
- $\langle 1 \rangle 6$ . Pick c such that, for all  $x \in A$ ,  $x \cap c = \emptyset$
- $\langle 1 \rangle 7. \ c \cap \bigcup A \in V_{\alpha}$
- $\langle 1 \rangle 8. \ \forall x \in V_{\alpha}(x \in A \Rightarrow \exists y \in V_{\alpha} \forall z \in V_{\alpha}(z \in x \land z \in c \cap \bigcup A \Leftrightarrow z = y))$

**Theorem 11.0.10.** For any ordinal  $\alpha$ , we have  $V_{\alpha}$  is a model of the Axiom of Regularity.

### Proof:

- $\langle 1 \rangle 1$ . Let:  $\alpha$  be an ordinal.
- $\langle 1 \rangle 2$ . Let:  $A \in V_{\alpha}$
- $\langle 1 \rangle 3$ . Assume:  $\exists x \in V_{\alpha}.x \in A$
- $\langle 1 \rangle 4$ . Pick  $m \in A$  of least rank.
- $\langle 1 \rangle 5. \ m \in V_{\alpha}$
- $\langle 1 \rangle 6. \ \neg \exists x \in V_{\alpha} (x \in m \land x \in A)$

**Theorem Schema 11.0.11.** For any axiom  $\alpha$  of Zermelo set theory, the following is a theorem:

For any limit ordinal  $\lambda > \omega$ , we have  $V_{\lambda}$  is a model of  $\alpha$ .

PROOF: Theorems 11.0.2, 11.0.3, 11.0.4, 11.0.5, 11.0.6, 11.0.7, 11.0.8, 11.0.9, 11.0.10.  $\Box$ 

Corollary Schema 11.0.11.1. for any axiom  $\alpha$  of Zermelo set theory, the following is a theorem:

 $V_{\omega 2}$  is a model of  $\alpha$ .

# Group Theory

## 12.1 Groups

**Definition 12.1.1** (Group). A group G consists of a set G and a function  $\cdot: G^2 \to G$  such that:

- $1. \cdot is associative$
- 2. There exists  $e \in G$  such that  $\forall x \in G.xe = x$  and  $\forall x \in G.\exists y \in G.xy = e$ .

Proposition 12.1.2. The inverse of an element in a group is unique.

Proof:

 $\langle 1 \rangle 1$ . Assume: b and b' are inverses of a.

 $\langle 1 \rangle 2. \ b = b'$ 

Proof:

$$b = be$$

$$= bab'$$

$$= eb'$$

$$= b'$$

**Definition 12.1.3.** We write  $x^{-1}$  for the inverse of x.

**Proposition 12.1.4.** In any group, if ab = ac then b = c.

Proof:

$$b = eb$$

$$= a^{-1}ab$$

$$= a^{-1}ac$$

$$= ec$$

$$= c$$

## 12.2 Abelian Groups

**Definition 12.2.1** (Abelian group). An  $Abelian\ group$  is a group whose multiplication is commutative.

We may say we are writing an Abelian group *additively*, meaning we write a + b for ab, 0 for e and -a for  $a^{-1}$ . In this case we write a - b for  $ab^{-1}$ .

# Ring Theory

### 13.1 Rings

**Definition 13.1.1** (Commutative Ring). A *commutative ring* consists of a set R and two binary operations +,  $\cdot$  on R such that:

- D is an Abelian group under +. Let us write 0 for its identity element.
- $\bullet$  · is commutative and associative, and distributes over +.
- $\bullet$  · has an identity element 1 that is different from 0.

**Proposition 13.1.2.** In any commutative ring, 0x = 0.

Proof:

$$(0+0)x = 0x$$

$$\therefore 0x + 0x = 0x + 0$$

$$\therefore 0x = 0 \qquad \text{(Proposition 12.1.4)} \square$$

**Proposition 13.1.3.** In any commutative ring, (-a)b = -(ab).

Proof:

$$ab + (-a)b = (a + (-a))b$$
  
=  $0b$   
=  $0$  (Proposition 13.1.2) $\square$ 

## 13.2 Ordered Rings

**Definition 13.2.1** (Ordered Commutative Ring). An ordered commutative ring consists of a commutative ring R with a linear order < on R such that:

• for all  $x, y, z \in R$ , we have x < y if and only if x + z < y + z.

• for all  $x, y, z \in R$ , if 0 < z then we have x < y if and only if xz < yz.

**Proposition 13.2.2.** In any ordered commutative ring, 0 < 1.

PROOF: If 1 < 0 then we have 0 < -1 and so 0 < (-1)(-1) = 1, which is a contradiction.  $\square$ 

**Proposition 13.2.3.** The ordering on an ordered commutative ring is dense; that is, if x < y then there exists z such that x < z < y.

PROOF: Take z = (x + y)/2.  $\square$ 

## 13.3 Integral Domains

**Definition 13.3.1** (Integral Domain). An *integral domain* is a commutative ring such that, for all  $a, b \in D$ , if ab = 0 then a = 0 or b = 0.

**Proposition 13.3.2.** In any integral domain, if ab = ac and  $a \neq 0$  then b = c.

PROOF: We have a(b-c)=0 and  $a\neq 0$  so b-c=0 hence b=c.  $\square$ 

**Definition 13.3.3** (Ordered Integral Domain). An *ordered integral domain* is an ordered commutative ring that is an integral domain.

# Field Theory

### 14.1 Fields

**Definition 14.1.1** (Field). A *field* F is a commutative ring such that  $0 \neq 1$  and, for all  $x \in F$ , if  $x \neq 0$  then there exists  $y \in F$  such that xy = 1.

Proposition 14.1.2. Every field is an integral domain.

PROOF: If ab = 0 and  $a \neq 0$  then  $b = a^{-1}ab = 0$ .  $\square$ 

**Proposition 14.1.3.** In any field F, we have  $F - \{0\}$  is an Abelian group under multiplication.

PROOF: Immediate from the definition.  $\Box$ 

**Definition 14.1.4** (Field of Fractions). Let D be an integral domain. The *field* of fractions of D is the quotient set  $F = (D \times (D - \{0\})) / \sim$  where

$$(a,b) \sim (c,d) \Leftrightarrow ad = bc$$

under

$$[(a,b)] + [(c,d)] = [(ad+bc,bd)]$$
$$[(a,b)][(c,d)] = [(ac,bd)]$$

We prove this is a field.

Proof:

 $\langle 1 \rangle 1$ .  $\sim$  is an equivalence relation on  $D \times (D - \{0\})$ . PROOF:

 $\langle 2 \rangle 1. \sim \text{ is reflexive.}$ 

PROOF: We always have ab = ba.

 $\langle 2 \rangle 2$ .  $\sim$  is symmetric.

PROOF: If ad = bc then cb = da.

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\langle 2 \rangle 3. \sim is transitive.
     \langle 3 \rangle 1. Assume: (a,b) \sim (c,d) \sim (e,f)
     \langle 3 \rangle 2. ad = bc and cf = de
     \langle 3 \rangle 3. adf = bde
        PROOF: adf = bcf = bde
     \langle 3 \rangle 4. af = be
        Proof: Proposition 13.3.2.
\langle 1 \rangle 2. Addition is well-defined.
  Proof:
  \langle 2 \rangle 1. If b \neq 0 and d \neq 0 then bd \neq 0.
     PROOF: Since D is an integral domain.
  \langle 2 \rangle 2. If ab' = a'b and cd' = c'd then (ad + bc)b'd' = (a'd' + b'c')bd.
     Proof:
                                 (ad + bc)b'd' = ab'dd' + bb'cd'
                                                 = a'bdd' + bb'c'd
                                                 = (a'd' + b'c')bd
\langle 1 \rangle 3. Multiplication is well-defined.
  Proof:
  \langle 2 \rangle 1. If b \neq 0 and d \neq 0 then bd \neq 0.
     PROOF: Since D is an integral domain.
  (2)2. If [(a,b)] = [(a',b')] and [(c,d)] = [(c',d')] then [(ac,bd)] = [(a'c',b'd')].
     PROOF: If ab' = a'b and cd' = c'd then acb'd' = a'c'bd.
\langle 1 \rangle 4. Addition is commutative.
  PROOF: [(a,b)] + [(c,d)] = [(ad+bc,bd)] = [(cb+da,db)] = [(c,d)] + [(a,b)] \sqcup
\langle 1 \rangle5. Addition is associative.
  Proof:
          [(a,b)] + ([(c,d)] + [(e,f)]) = [(a,b)] + [(cf + de, df)]
                                            = [(adf + bcf + bde, bdf)]
                                            = [(ad + bc, bd)] + [(e, f)]
                                            = ([(a,b)] + [(c,d)]) + [(e,f)]
                                                                                       \langle 1 \rangle 6. For any x \in F we have x + [(0,1)] = x
  PROOF: [(a,b)] + [(0,1)] = [(a \cdot 1 + b \cdot 0, b \cdot 1)] = [(a,b)] \sqcup
\langle 1 \rangle 7. For any x \in F, there exists y \in F such that x + y = [(0,1)].
  PROOF: [(a,b)] + [(-a,b)] = [(ab-ab,b^2)] = [(0,b^2)] = [(0,1)]
\langle 1 \rangle 8. Multiplication is commutative.
  PROOF: [(a,b)][(c,d)] = [(c,d)][(a,b)] = [(ac,bd)].
\langle 1 \rangle 9. Multiplication is assocative.
  PROOF: [(a,b)]([(c,d)][(e,f)]) = ([(a,b)][(c,d)])[(e,f)] = [(ace,bdf)].
\langle 1 \rangle 10. For any x \in F we have x[(1,1)] = x
  PROOF: [(a,b)][(1,1)] = [(a,b)]
\langle 1 \rangle 11. For any non-zero x \in F, there exists y \in F such that xy = [(1,1)].
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Proof:
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 \begin{array}{l} \langle 2 \rangle 1. \ \ \mathrm{Let:} \ \left[ (a,b) \right] \in \mathbb{Q} \\ \langle 2 \rangle 2. \ \ \mathrm{Assume:} \ \left[ (a,b) \right] \neq \left[ (0,1) \right] \\ \langle 2 \rangle 3. \ \ a \neq 0 \\ \langle 2 \rangle 4. \ \ \left[ (a,b) \right] \left[ (b,a) \right] = \left[ (1,1) \right] \\ \square \\ \end{array}
```

**Definition 14.1.5.** For any field F, let N(F) be the intersection of all the subsets  $S \subseteq F$  such that  $1 \in S$  and  $\forall x \in S.x + 1 \in S$ .

**Definition 14.1.6** (Characteristic Zero). A field F has *characteristic* 0 iff  $0 \notin N(F)$ .

**Proposition 14.1.7.** In a field F with characteristic 0, the function  $n : \mathbb{N} \to N(F)$  defined by

$$n(0) = 1$$
$$n(x+1) = n(x) + 1$$

is a bijection.

### Proof:

 $\langle 1 \rangle 1$ . *n* is injective.

 $\langle 2 \rangle 1$ . Assume: for a contradiction n(i) = n(j) with  $i \neq j$ 

 $\langle 2 \rangle 2$ . Assume: w.l.o.g. i < j

 $\langle 2 \rangle 3$ . n(j-i)=0

 $\langle 2 \rangle 4$ . Q.E.D.

PROOF: This contradicts the fact that F has characteristic 0.

 $\langle 1 \rangle 2$ . *n* is surjective.

PROOF: Since ran n is a subset of F that includes 1 and is closed under +1.

**Definition 14.1.8.** In any field F, let

$$I(F) = N(F) \cup \{0\} \cup \{-x \mid x \in N(F)\}\$$

**Definition 14.1.9.** In any field F, let

$$Q(F) = \{x/y \mid x, y \in I(F), y \neq 0\}$$

**Proposition 14.1.10.** Q(F) is the smallest subfield of F.

PROOF: Q(F) is closed under + and  $\cdot$ , and any subset of F closed under + and  $\cdot$  that contains 0 and 1 must include Q(F).  $\square$ 

**Theorem 14.1.11.** Let F and G be fields of characteristic 0. Then there exists a unique field isomorphism between Q(F) and Q(G).

Proof:

- $\langle 1 \rangle 1$ . Let:  $\phi: N(F) \to N(G)$  be the unique function such that  $\phi(1) = 1$  and  $\forall x \in N(F). \phi(x+1) = \phi(x) + 1$ .
- $\langle 1 \rangle 2$ .  $\phi$  is a bijection.

Proof: Similar to Proposition 14.1.7.

 $\langle 1 \rangle 3. \ \forall x, y \in N(F). \phi(x+y) = \phi(x) + \phi(y)$ 

Proof: Induction on y.

 $\langle 1 \rangle 4. \ \forall x, y \in N(F). \phi(xy) = \phi(x)\phi(y)$ 

PROOF: Induction on y.

- $\langle 1 \rangle$ 5. Extend  $\phi$  to a bijection  $I(F) \cong I(G)$  such that  $\forall x, y \in I(F).\phi(x+y) = \phi(x) + \phi(y)$  and  $\forall x, y \in I(F).\phi(xy) = \phi(x)\phi(y)$ 
  - $\langle 2 \rangle 1$ . Define  $\phi(0) = 0$  and  $\phi(-x) = -\phi(x)$  for  $x \in N(F)$ 
    - $\langle 3 \rangle 1. \ 0 \notin N(F)$
    - $\langle 3 \rangle 2$ . For all  $x \in N(F)$  we have  $-x \notin N(F)$

PROOF: Then we would have  $x + -x = 0 \in N(F)$ .

- $\langle 3 \rangle 3$ . For all  $x \in N(F)$  we have  $-x \neq 0$
- $\langle 2 \rangle 2$ . For all  $x, y \in I(F)$  we have  $\phi(x+y) = \phi(x) + \phi(y)$

PROOF: Case analysis on x and y.

 $\langle 2 \rangle 3$ . For all  $x, y \in I(F)$  we have  $\phi(xy) = \phi(x)\phi(y)$ 

PROOF: Case analysis on x and y.

- $\langle 1 \rangle$ 6. Extend  $\phi$  to a bijection  $Q(F) \cong Q(G)$  such that  $\forall x, y \in Q(F).\phi(x+y) = \phi(x) + \phi(y)$  and  $\forall x, y \in Q(F).\phi(xy) = \phi(x)\phi(y)$ 
  - $\langle 2 \rangle 1$ . Define  $\phi(x/y) = \phi(x)/\phi(y)$
- $\langle 1 \rangle 7$ .  $\phi$  is unique.
  - $\langle 2 \rangle 1$ . Let:  $\theta$  satisfy the theorem.
  - $\langle 2 \rangle 2$ . For all  $x \in N(F)$  we have  $\theta(x) = \phi(x)$
  - $\langle 2 \rangle 3$ . For all  $x \in I(F)$  we have  $\theta(x) = \phi(x)$
  - $\langle 2 \rangle 4$ . For all  $x \in Q(F)$  we have  $\theta(x) = \phi(x)$

### 14.2 Ordered Fields

**Definition 14.2.1** (Ordered Field). An *ordered field* is an ordered commutative ring that is a field.

**Proposition 14.2.2.** Every ordered field F has characteristic  $\theta$ .

PROOF: We have 0 < n for all  $n \in N(F)$ .  $\square$ 

**Proposition 14.2.3.** Let F be a field of characteristic 0. Then there exists a unique relation < on Q(F) that makes Q(F) into an ordered field.

Proof: Easy.  $\square$ 

Corollary 14.2.3.1. Let F and G be ordered fields. Let  $\phi$  be the unique field isomorphism between Q(F) and Q(G). Then  $\phi$  is an ordered field isomorphism.

**Definition 14.2.4** (Archimedean). An ordered field F is Archimedean iff

$$\forall x \in F. \exists n \in N(F). n > x .$$

**Proposition 14.2.5.** Let F be an Archimedean ordered field. Let  $x, y \in F$  with x > 0. Then there exists  $n \in N(F)$  such that nx > y.

PROOF: Pick n > y/x.  $\square$ 

**Proposition 14.2.6.** Let F be an Archimedean ordered field. For all  $x, y \in F$ , if x < y, then there exists  $r \in Q(F)$  such that x < r < y.

### Proof:

- $\langle 1 \rangle 1$ . Case: x > 0
  - $\langle 2 \rangle 1$ . PICK  $n \in N(F)$  such that n(y-x) > 1

Proof: Proposition 14.2.5.

- $\langle 2 \rangle 2$ . ny > 1 + nx
- $\langle 2 \rangle$ 3. Let: m be the least element of N(F) such that m > nx.
- $\langle 2 \rangle 4$ .  $m-1 \leq nx$
- $\langle 2 \rangle 5$ . nx < m < ny
- $\langle 2 \rangle 6$ . x < m/n < y
- $\langle 1 \rangle 2$ . Case:  $x \leq 0$ 
  - $\langle 2 \rangle 1$ . PICK  $k \in N(F)$  such that k > -x
  - $\langle 2 \rangle 2$ . 0 < x + k < y + k
  - $\langle 2 \rangle 3$ . PICK  $r \in Q(F)$  such that x + k < r < y + k

Proof:  $\langle 1 \rangle 1$ 

 $\langle 2 \rangle 4$ . x < r - k < y

**Definition 14.2.7** (Complete). An ordered field F is *complete* iff every nonempty subset of F bounded above has a least upper bound.

Proposition 14.2.8. Every complete ordered field is Archimedean.

#### Proof:

- $\langle 1 \rangle 1$ . Let: F be a complete ordered field.
- $\langle 1 \rangle 2$ . Let:  $x \in F$
- $\langle 1 \rangle$ 3. Assume: for a contradiction there is no member of N(F) greater than x.
- $\langle 1 \rangle 4$ . x is an upper bound for N(F).
- $\langle 1 \rangle 5$ . Let:  $y = \sup N(F)$
- $\langle 1 \rangle 6$ . Pick  $n \in N(F)$  such that y 1 < n
- $\langle 1 \rangle 7$ . y < n+1
- $\langle 1 \rangle 8$ . Q.E.D.

Proof: This is a contradiction.

**Proposition 14.2.9.** Let F be a complete ordered field and  $a \in F$  be nonnegative. Then there exists  $b \in F$  such that  $b^2 = a$ .

### Proof:

- $\langle 1 \rangle 1$ . Let:  $B = \{ x \in F \mid 0 \le x \le 1 + a \}$
- $\langle 1 \rangle 2$ . Let:  $\phi : B \to B$  be the function

$$\phi(x) = x + \frac{1}{2(1+a)}(a-x^2) .$$

- $\langle 1 \rangle 3$ .  $\phi$  is strictly monotone.
  - $\langle 2 \rangle$ 1. Let:  $0 \le x < y \le 1 + a$  $\langle 2 \rangle$ 2.  $1 \frac{x+y}{2(1+a)} > 0$

  - $\langle 2 \rangle 3. \ \phi(y) \phi(x) = (y x)(1 \frac{x+y}{2(1+a)}) > 0$
  - $\langle 2 \rangle 4. \ \phi(x) < \phi(y)$
- $\langle 1 \rangle 4$ . Pick  $b \in B$  such that  $\phi(b) = b$ .

PROOF: Knaster Fixed-Point Theorem.

$$\langle 1 \rangle 5. \ b^2 = a$$

**Theorem 14.2.10** (Uniqueness of the Complete Ordered Field). If F and G are complete ordered fields, then there exists a unique bijection  $\phi: F \cong G$  such that, for all  $x, y \in F$ ,

$$\phi(x+y) = \phi(x) + \phi(y)$$
$$\phi(xy) = \phi(x)\phi(y)$$

This bijection also satisfies: for all  $x, y \in F$ ,

$$x < y \Leftrightarrow \phi(x) < \phi(y)$$
.

#### Proof:

 $\langle 1 \rangle 1$ . Pick a bijection  $\phi: Q(F) \cong Q(G)$  such that, for all  $x, y \in Q(F)$ ,

$$\phi(x+y) = \phi(x) + \phi(y)$$
$$\phi(xy) = \phi(x)\phi(y)$$

$$x < y \Leftrightarrow \phi(x) < \phi(y)$$

Proof: Corollary 14.2.3.1.

 $\langle 1 \rangle 2$ . Q(F) intersects every interval in F.

Proof: Proposition 14.2.6.

 $\langle 1 \rangle 3$ . Q(G) intersects every interval in G.

Proof: Proposition 14.2.6.

 $\langle 1 \rangle 4$ . PICK an order isomorphism  $\psi : F \cong G$  that extends  $\phi$ .

PROOF: Theorem 5.1.21.

- $\langle 1 \rangle 5. \ \forall x, y \in F. \psi(x+y) = \psi(x) + \psi(y)$ 
  - $\langle 2 \rangle 1$ . Let:  $x, y \in F$
  - $\langle 2 \rangle 2$ .  $\psi(x) + \psi(y) \not< \psi(x+y)$ 
    - $\langle 3 \rangle 1$ . Assume: for a contradiction  $\psi(x) + \psi(y) < \psi(x+y)$
    - $\langle 3 \rangle 2$ . Pick  $r' \in Q(G)$  such that  $\psi(x) < r' < \psi(x+y) \psi(y)$
    - $\langle 3 \rangle 3$ . Pick  $s' \in Q(G)$  such that  $\psi(y) < s' < \psi(x+y) r'$
    - $\langle 3 \rangle 4. \ r' + s' < \psi(x+y)$
    - $\langle 3 \rangle 5$ . Pick  $r, s \in Q(F)$  such that  $\phi(r) = r'$  and  $\phi(s) = s'$
    - $\langle 3 \rangle 6. \ \phi(r+s) = r' + s'$
    - $\langle 3 \rangle 7. \ \psi(x) < \psi(r)$
    - $\langle 3 \rangle 8. \ \psi(y) < \psi(s)$
    - $\langle 3 \rangle 9. \ \psi(x+y) > \psi(r+s)$
    - $\langle 3 \rangle 10. \ x < r$

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\langle 3 \rangle 11. \ y < s
       \langle 3 \rangle 12. x + y > r + s
       \langle 3 \rangle 13. Q.E.D.
           PROOF: This is a contradiction.
   \langle 2 \rangle 3. \ \psi(x+y) \not< \psi(x) + \psi(y)
       Proof: Similar.
\langle 1 \rangle 6. \ \forall x, y \in F. \psi(xy) = \psi(x) \psi(y)
    \langle 2 \rangle 1. Let: x, y \in F
   \langle 2 \rangle 2. Case: x and y are positive.
       \langle 3 \rangle 1. \ \psi(x)\psi(y) \not< \psi(xy)
           \langle 4 \rangle1. Assume: for a contradiction \psi(x)\psi(y) < \psi(xy)
           \langle 4 \rangle 2. PICK r' \in Q(G) such that \psi(x) < r' < \psi(xy)/\psi(y)
           \langle 4 \rangle 3. Pick s' \in Q(G) such that \psi(y) < s' < \psi(xy)/r'
           \langle 4 \rangle 4. r's' < \psi(xy)
           \langle 4 \rangle5. PICK r, s \in Q(F) such that \phi(r) = r' and \phi(s) = s'
           \langle 4 \rangle 6. \ \phi(rs) = r's'
           \langle 4 \rangle 7. x < r, y < s \text{ and } rs < xy
           \langle 4 \rangle 8. Q.E.D.
              PROOF: This is a contradiction.
       \langle 3 \rangle 2. \ \psi(xy) \not< \psi(x)\psi(y)
           PROOF: Similar.
   \langle 2 \rangle 3. Case: x and y are not both positive.
       PROOF: Follows from \langle 2 \rangle 2 since \psi(-x) = -\psi(x) by \langle 1 \rangle 5.
\langle 1 \rangle 7. For any field isomorphism \theta : F \cong G, we have \theta = \psi.
   \langle 2 \rangle 1. \ \theta \upharpoonright Q(F) = \phi
       PROOF: Theorem 14.1.11.
   \langle 2 \rangle 2. \theta is strictly monotone.
       \langle 3 \rangle 1. Let: x, y \in F with x < y
       \langle 3 \rangle 2. y - x > 0
       \langle 3 \rangle 3. Pick z \in F such that z^2 = y - x
       \langle 3 \rangle 4. \theta(z)^2 = \theta(y) - \theta(x)
       \langle 3 \rangle 5. \theta(y) - \theta(x) > 0
       \langle 3 \rangle 6. \ \theta(x) < \theta(y)
   \langle 2 \rangle 3. \ \theta = \psi
       Proof: By the uniqueness of \psi.
```

# Chapter 15

# Number Systems

## 15.1 The Integers

**Definition 15.1.1.** The set of integers  $\mathbb{Z}$  is the quotient set  $\mathbb{N}^2/\sim$ , where  $(m,n)\sim(p,q)$  iff m+q=n+p.

We prove  $\sim$  is an equivalence relation on  $\mathbb{N}^2$ .

#### Proof:

 $\langle 1 \rangle 1$ .  $\sim$  is reflexive.

PROOF: For all  $m, n \in \mathbb{N}$  we have m + n = n + m.

 $\langle 1 \rangle 2$ .  $\sim$  is symmetric.

PROOF: If m + q = n + p then p + n = q + m.

- $\langle 1 \rangle 3$ .  $\sim$  is transitive.
  - $\langle 2 \rangle 1$ . Assume:  $(m,n) \sim (p,q) \sim (r,s)$
  - $\langle 2 \rangle 2$ . m+q=n+p and p+s=q+r
  - $\langle 2 \rangle 3$ . m+q+s=n+q+r
  - $\langle 2 \rangle 4$ . m+s=n+r

PROOF: By cancellation.

**Definition 15.1.2** (Addition). Define  $addition + \text{ on } \mathbb{Z}$  by [(m,n)] + [(p,q)] = [(m+p,n+q)].

We prove this is well-defined.

PROOF: If m+n'=n+m' and p+q'=q+p' then m+p+n'+q'=n+q+m'+p'.

**Proposition 15.1.3.** Addition on  $\mathbb{Z}$  is commutative.

PROOF: 
$$[(m,n)] + [(p,q)] = [(m+p,n+q)] = [(p+m,q+n)] = [(p,q)] + [(m,n)].$$

**Proposition 15.1.4.** Addition on  $\mathbb{Z}$  is associative.

PROOF: [(m,n)] + ([(p,q)] + [(r,s)]) = [(m+p+r,n+q+s)] = ([(m,n)] + [(p,q)]) + [(r,s)].

**Proposition 15.1.5.** Given natural numbers m and n, we have [(m,0)] = [(n,0)] iff m = n.

PROOF: Immediate from definitions.

**Definition 15.1.6.** We identify any natural number n with the integer [(n,0)].

**Proposition 15.1.7.** Addition on integers agrees with addition on natural numbers.

PROOF: Since [(m,0)] + [(n,0)] = [(m+n,0)].

**Proposition 15.1.8.** For all  $a \in \mathbb{Z}$  we have a + 0 = a.

PROOF: [(m, n)] + [(0, 0)] = [(m + 0, n + 0)] = [(m, n)].

**Proposition 15.1.9.** For all  $a \in \mathbb{Z}$ , there exists  $b \in \mathbb{Z}$  such that a + b = 0.

PROOF: [(m,n)] + [(n,m)] = [(m+n,m+n)] = [(0,0)]

Proposition 15.1.10. The integers form an Abelian group under addition.

PROOF: Proposition 15.1.3, 15.1.4, 15.1.8, 15.1.9.

**Definition 15.1.11.** Define multiplication  $\cdot$  on  $\mathbb{Z}$  by: [(m,n)][(p,q)] = [(mp + nq, mq + np)].

We prove this is well defined.

Proof:

- $\langle 1 \rangle 1.$  Assume: m+n'=n+m' and p+q'=q+p' Prove: mp+nq+m'q'+n'p'=mq+np+m'p'+n'q'
- $\langle 1 \rangle 2$ . mp + n'p = np + m'p
- $\langle 1 \rangle 3$ . nq + m'q = mq + n'q
- $\langle 1 \rangle 4. \ m'p + m'q' = m'q + m'p'$
- $\langle 1 \rangle 5. \ n'q + n'p' = n'p + n'q'$
- (1)6. mp + n'p + nq + m'q + m'p + m'q' + n'q + n'p' = np + m'p + mq + n'q + m'q + m'p' + n'p + n'q'
- $\langle 1 \rangle 7$ . mp + nq + m'q' + n'p' = mq + np + m'p' + n'q'

Proof: By cancellation.

**Proposition 15.1.12.** Multiplication on integers agrees with multiplication on natural numbers.

PROOF: Since [(m,0)][(n,0)] = [(mn+0,m0+n0)] = [(mn,0)].

**Proposition 15.1.13.** *Multiplication on*  $\mathbb{Z}$  *is commutative.* 

Proof: [(m,n)][(p,q)] = [(mp+nq,mq+np)] = [(pm+qn,pn+qm)] = [(p,q)][(m,n)].

**Proposition 15.1.14.** *Multiplication on*  $\mathbb{Z}$  *is associative.* 

Proof:

$$\begin{split} [(m,n)]([(p,q)][(r,s)]) &= [(m,n)][(pr+qs,ps+qr)] \\ &= [(mpr+mqs+nps+nqr,mps+mqr+npr+nqs)] \\ &= [(mp+nq,mq+np)][(r,s)] \\ &= ([(m,n)][(p,q)])[(r,s)] \end{split}$$

Proposition 15.1.15. Multiplication distributes over addition.

Proof:

$$\begin{split} [(m,n)]([(p,q)]+[(r,s)]) &= [(m,n)][(p+r,q+s)] \\ &= [(mp+mr+nq+ns,np+nr+mq+ms)] \\ [(m,n)][(p,q)]+[(m,n)][(r,s)] &= [(mp+nq,mq+np)]+[(mr+ns,ms+nr)] \\ &= [(mp+nq+mr+ns,mq+np+ms+nr)] \end{split}$$

**Proposition 15.1.16.** For any integer a we have a1 = a.

PROOF: Since 
$$[(m,n)][(1,0)] = [(m1+n0,m0+n1)] = [(m,n)]$$
.

**Proposition 15.1.17.** For any integers a and b, if ab = 0 then a = 0 or b = 0.

Proof:

```
\langle 1 \rangle 1. Assume: [(m,n)][(p,q)] = [(0,0)]
\langle 1 \rangle 2. mp + nq = mq + np
\langle 1 \rangle 3. Assume: [(m,n)] \neq [(0,0)]
\langle 1 \rangle 4. \ m \neq n
        Prove: p = q
\langle 1 \rangle 5. Case: m < n
   \langle 2 \rangle 1. \ p \not < q
      PROOF: If p < q then mq + np < mp + nq by Proposition 8.4.6.
```

PROOF: If q < p then mp + nq < mq + np by Proposition 8.4.6.

 $\langle 2 \rangle 3. \ p = q$ 

PROOF: By trichotomy.

 $\langle 1 \rangle 6$ . Case: n < m

PROOF: Similar.

**Proposition 15.1.18.** The integers  $\mathbb{Z}$  form an integral domain.

PROOF: Propositions 15.1.13, 15.1.14, 15.1.15, 15.1.16, 15.1.17, 15.1.10.

**Definition 15.1.19.** Define < on  $\mathbb{Z}$  by [(m,n)] < [(p,q)] if and only if m+q <n+p.

We prove this is well-defined.

$$\begin{array}{ll} \langle 1 \rangle 1. & \text{Assume: } m+n'=n+m' \text{ and } p+q'=q+p'. \\ & \text{Prove: } m+q< n+p \text{ if and only if } m'+q'< n'+p' \\ \langle 1 \rangle 2. & m+q< n+p \text{ if and only if } m'+q'< n'+p' \\ & \text{Proof: } \\ & m+q< n+p \Leftrightarrow m+n'+q< n+n'+p \\ & \Leftrightarrow m'+n+q< n+n'+p \\ & \Leftrightarrow m'+q< n'+p \\ & \Leftrightarrow m'+q+p'< n'+p+p' \end{array} \qquad \begin{array}{l} \text{(Corollary 6.5.7.1)} \\ \text{(Corollary 6.5.7.1)} \\ & \Leftrightarrow m'+q'+p+p' \end{array}$$

**Proposition 15.1.20.** The ordering on the integers agrees with the ordering on the natural numbers.

PROOF: We have [(m,0)] < [(n,0)] iff m < n.  $\square$ 

**Proposition 15.1.21.** < is a linear order on  $\mathbb{Z}$ .

#### Proof:

 $\langle 1 \rangle 1$ . < is irreflexive.

PROOF: We never have m + n < m + n.

- $\langle 1 \rangle 2$ . < is transitive.
  - $\langle 2 \rangle 1$ . Assume: [(m,n)] < [(p,q)] < [(r,s)]
  - $\langle 2 \rangle 2$ . m+q < n+p and p+s < q+r
  - $\langle 2 \rangle 3. \ m + q + s < n + q + r$

PROOF: m + q + s < n + p + s < n + q + r

 $\langle 2 \rangle 4$ . m+s < n+r

PROOF: Corollary 6.5.7.1.

 $\langle 1 \rangle 3.$  < is total.

PROOF: Given natural numbers m, n, p and q, either m+q < n+p, or m+q=n+p, or n+p < m+q.

**Definition 15.1.22** (Positive). An integer a is positive iff a > 0.

**Theorem 15.1.23.** For any integers a, b and c, we have a < b if and only if a + c < b + c.

- $\langle 1 \rangle 1$ . If a < b then a + c < b + c.
  - $\langle 2 \rangle 1$ . Let: a = [(m, n)], b = [(p, q)] and c = [(r, s)].
  - $\langle 2 \rangle 2$ . Assume: a < b
  - $\langle 2 \rangle 3. \ m+q < n+p$
  - $\langle 2 \rangle 4$ . m + r + q + s < n + r + p + s
  - $\langle 2 \rangle 5. [(m+r, n+s)] < [(p+r, q+s)]$
  - $\langle 2 \rangle 6$ . a+c < b+c

```
\langle 1 \rangle2. If a+c < b+c then a < b.
PROOF: From \langle 1 \rangle1 and Proposition 5.2.6.
```

**Proposition 15.1.24.** Let a, b and c be integers. If 0 < c, then a < b if and only if ac < bc.

```
PROOF: \langle 1 \rangle 1. Let: c = [(r, s)] \langle 1 \rangle 2. Assume: 0 < c \langle 1 \rangle 3. s < r \langle 1 \rangle 4. For all integers a and b, if a < b then ac < bc \langle 2 \rangle 1. Let: a = [(m, n)], b = [(p, q)]. \langle 2 \rangle 2. Assume: a < b \langle 2 \rangle 3. m + q < n + p \langle 2 \rangle 4. (m + q)r + (p + n)s < (m + q)s + (p + n)r Proof: Proposition 8.4.6, \langle 1 \rangle 3, \langle 2 \rangle 3. \langle 2 \rangle 5. mr + ns + ps + qr < ms + nr + pr + qs \langle 2 \rangle 6. [(mr + ns, ms + nr)] < [(pr + qs, ps + qr)] \langle 2 \rangle 7. ac < bc \langle 1 \rangle 5. For all integers a and b, if ac < bc then a < b Proof: From \langle 1 \rangle 4 and Proposition 5.2.6.
```

**Proposition 15.1.25.** Let a be a positive integer. For any integer b, there exists  $k \in \mathbb{N}$  such that b < ak.

```
PROOF: \langle 1 \rangle 1. Case: b \leq 0
PROOF: Take k = 1. \langle 1 \rangle 2. Case: b > 0
PROOF: Take k = b + 1.
```

#### 15.2 The Rationals

**Definition 15.2.1** (Rational Numbers). The set  $\mathbb{Q}$  of rational numbers is the field of fractions over the integers.

**Proposition 15.2.2.** For any integers a and b, we have [(a,1)] = [(b,1)] iff a = b.

Proof: Immediate from definitions.

Henceforth we identify any integer a with the rational number [(a, 1)].

**Proposition 15.2.3.** Addition on the rationals agrees with addition on the integers.

PROOF: 
$$[(a, 1)] + [(b, 1)] = [(a \cdot 1 + b \cdot 1, 1 \cdot 1)] = [(a + b, 1)].$$

**Proposition 15.2.4.** Multiplication on the rationals agrees with multiplication on the integers.

PROOF: 
$$[(a, 1)][(b, 1)] = [(ab, 1)]$$

**Definition 15.2.5.** Define the ordering < on the rationals by: if b and d are positive, then [(a,b)] < [(c,d)] iff ad < bc.

We prove this is well-defined.

#### Proof:

 $\langle 1 \rangle 1$ . For any rational q, there exist integers a, b with b positive such that q = [(a, b)].

PROOF: Since [(a,b)] = [(-a,-b)], and if  $b \neq 0$  then one of b and -b is positive.

 $\langle 1 \rangle 2$ . If b, b', d and d' are positive, [(a,b)] = [(a',b')], and [(c,d)] = [(c',d')], then ad < bc iff a'd' < b'c'.

#### Proof:

- $\langle 2 \rangle 1$ . If ad < bc then a'd' < b'c'.
  - $\langle 3 \rangle 1$ . Assume: ad < bc
  - $\langle 3 \rangle 2$ . ab'd < bb'c
  - $\langle 3 \rangle 3$ . a'bd < bb'c
  - $\langle 3 \rangle 4$ . a'd < b'c
  - $\langle 3 \rangle 5$ . a'dd' < b'cd'
  - $\langle 3 \rangle 6$ . a'dd' < b'c'd
  - $\langle 3 \rangle 7$ . a'd' < b'c'
- $\langle 2 \rangle 2$ . If a'd' < b'c' then ad < bc.

PROOF: Similar.

П

**Proposition 15.2.6.** The ordering on the rationals agrees with the ordering on the integers.

PROOF: We have [(a,1)] < [(b,1)] if and only if a < b.  $\square$ 

**Proposition 15.2.7.** The relation < is a linear ordering on  $\mathbb{Q}$ .

#### Proof:

 $\langle 1 \rangle 1$ . < is irreflexive.

PROOF: We never have ab < ab.

- $\langle 1 \rangle 2$ . < is transitive.
  - $\langle 2 \rangle 1$ . Assume: [(a,b)] < [(c,d)] < [(e,f)] where b, d and f are positive.
  - $\langle 2 \rangle 2$ . ad < bc and cf < de
  - $\langle 2 \rangle 3$ . adf < bde

Proof: adf < bcf < bde

- $\langle 2 \rangle 4$ . af < be
- $\langle 1 \rangle 3. < \text{is total.}$

PROOF: For any integers a, b, c, d, we have ad < bc or ad = bc or bc < ad.

**Proposition 15.2.8.** For any rationals r, s and t, we have r < s if and only if r + t < s + t.

Proof:

 $\langle 1 \rangle 1$ . Let: a, b, c, d, e, f be integers with b, d and f positive.

 $\langle 1 \rangle 2$ . [(a,b)] + [(e,f)] < [(c,d)] + [(e,f)] if and only if [(a,b)] < [(c,d)]. PROOF:

$$\begin{split} [(a,b)] + [(e,f)] < [(c,d)] + [(e,f)] &\Leftrightarrow [(af+be,bf)] < [(cf+de,df)] \\ &\Leftrightarrow (af+be)df < (cf+de)bf \\ &\Leftrightarrow afdf + bedf < cfbf + debf \\ &\Leftrightarrow afdf < cfbf \\ &\Leftrightarrow ad < bc \\ &\Leftrightarrow [(a,b)] < [(c,d)] \end{split}$$

**Corollary 15.2.8.1.** For any rational r, we have r < 0 if and only if 0 < -r.

**Definition 15.2.9** (Absolute Value). For any rational r, the absolute value of r is defined by

$$|r| := \begin{cases} -r & \text{if } 0 < -r \\ r & \text{otherwise} \end{cases}$$

**Proposition 15.2.10.** For any rationals r, s and t, if t is positive then r < s iff rt < st.

Proof:

 $\langle 1 \rangle 1$ . Let: r = [(a,b)], s = [(c,d)] and t = [(e,f)] where b, d and f are positive.

 $\langle 1 \rangle 2$ . Assume: 0 < t

 $\langle 1 \rangle 3. \ e > 0$ 

 $\langle 1 \rangle 4$ . rt < st iff r < s

Proof:

$$rt < st \Leftrightarrow [(ae, bf)] < [(ce, df)]$$
  
 $\Leftrightarrow aedf < cebf$   
 $\Leftrightarrow ad < bc$   
 $\Leftrightarrow r < s$ 

Corollary 15.2.10.1. The rationals form an ordered field.

**Proposition 15.2.11.** *Let* p *be a positive rational. For any rational number* r, *there exists*  $k \in \mathbb{N}$  *such that* r < pk.

Proof:

 $\langle 1 \rangle 1$ . Let: p = a/b and r = c/d where a, b and d are positive.

```
\langle 1 \rangle2. PICK k \in \mathbb{N} such that bc < adk PROOF: Proposition 15.1.25. \langle 1 \rangle3. r < pk □
```

#### Proposition 15.2.12. $\mathbb{Q} \approx \mathbb{N}$

PROOF: Arrange the rationals in order 0/1, 1/1, 1/2, 0/2, -1/2, -1/1, -2/1, -2/2, -2/3, -1/3, 0/3, 1/3, 2/3, etc. then remove all duplicates.  $\Box$ 

#### 15.3 The Real Numbers

**Definition 15.3.1** (Cauchy Sequence). A Cauchy sequence is a sequence  $(q_n)$  of rationals such that, for every positive rational  $\epsilon$ , there exists  $k \in \mathbb{N}$  such that  $\forall m, n > k. |q_m - q_n| < \epsilon$ .

**Definition 15.3.2** (Dedekind Cut). A *Dedekind cut* is a set  $x \subseteq \mathbb{Q}$  such that:

- 1.  $\emptyset \neq x \neq \mathbb{Q}$
- 2. x is closed downwards.
- 3. x has no greatest member.

The set  $\mathbb{R}$  of *real numbers* is the set of Dedekind cuts.

**Proposition 15.3.3.** For any rational q, we have  $\{r \in \mathbb{Q} \mid r < q\} \in \mathbb{R}$ .

```
Proof:
```

- $\langle 1 \rangle 1$ . Let:  $q \in \mathbb{Q}$
- $\langle 1 \rangle 2$ . Let:  $q \downarrow = \{r \mid r < q\}$
- $\langle 1 \rangle 3. \ \ q \downarrow \neq \emptyset$

PROOF: We have  $q - 1 \in q \downarrow$ .

 $\langle 1 \rangle 4. \ \ q \downarrow \neq \mathbb{Q}$ 

PROOF: Since  $q \notin q \downarrow$ .

 $\langle 1 \rangle 5$ .  $q \downarrow$  is closed downwards.

PROOF: Trivial.

 $\langle 1 \rangle 6$ .  $q \downarrow$  has no greatest element.

PROOF: For all  $r \in q \downarrow$  we have  $r < (q+r)/2 \in q \downarrow$ .

**Proposition 15.3.4.** For rationals q and r, we have q = r if and only if  $\{s \in \mathbb{Q} \mid s < q\} = \{s \in \mathbb{Q} \mid s < r\}.$ 

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $q \downarrow = \{ s \in \mathbb{Q} \mid s < q \}$
- $\langle 1 \rangle 2$ . Let:  $r \downarrow = \{ s \in \mathbb{Q} \mid s < r \}$
- $\langle 1 \rangle 3$ . If q = r then  $q \downarrow = r \downarrow$

PROOF: Trivial.

```
\begin{split} &\langle 1 \rangle 4. \text{ If } q < r \text{ then } q \downarrow \neq r \downarrow \\ &\text{PROOF: We have } q \in r \downarrow \text{ and } q \notin q \downarrow. \\ &\langle 1 \rangle 5. \text{ If } r < q \text{ then } q \downarrow \neq r \downarrow \\ &\text{PROOF: We have } r \in q \downarrow \text{ and } q \notin q \downarrow. \\ &\square \end{split}
```

Henceforth we identify a rational q with the real number  $\{r \in \mathbb{Q} \mid r < q\}$ .

**Definition 15.3.5.** Define the ordering < on  $\mathbb{R}$  by: x < y iff  $x \subseteq y$ .

**Proposition 15.3.6.** The ordering on the reals agrees with the ordering on the rationals.

```
Proof:
\langle 1 \rangle 1. Let: q, r \in \mathbb{Q}
\langle 1 \rangle 2. Let: q \downarrow = \{ s \in \mathbb{Q} \mid s < q \}.
\langle 1 \rangle 3. Let: r \downarrow = \{ s \in \mathbb{Q} \mid s < r \}.
            Prove: q < r \text{ iff } q \downarrow \subsetneq r \downarrow
\langle 1 \rangle 4. If q < r then q \downarrow \subseteq r \downarrow
     \langle 2 \rangle 1. Assume: q < r
     \langle 2 \rangle 2. q \downarrow \subseteq r \downarrow
          Proof: If s < q then s < r.
     \langle 2 \rangle 3. \ \ q \downarrow \neq r \downarrow
          Proof: Proposition 15.3.4.
\langle 1 \rangle 5. If q \downarrow \subsetneq r \downarrow then q < r
     \langle 2 \rangle 1. Assume: q \downarrow \subsetneq r \downarrow
     \langle 2 \rangle 2. Pick s \in r \downarrow such that s \notin q \downarrow
     \langle 2 \rangle 3. \ q \leq s < r
```

**Proposition 15.3.7.** The ordering < is a linear ordering on  $\mathbb{R}$ .

```
Proof:
```

```
\langle 1 \rangle 1. < is irreflexive.
```

PROOF: No set is a proper subset of itself.

 $\langle 1 \rangle 2$ . < is transitive.

PROOF: Since the relationship  $\subseteq$  is transitive on the class of all sets.

- $\langle 1 \rangle 3$ . < is total.
  - $\langle 2 \rangle 1$ . Let: x, y be Dedekind cuts.
  - $\langle 2 \rangle 2$ . Assume:  $x \nsubseteq y$ Prove:  $y \subsetneq x$
  - $\langle 2 \rangle 3$ . PICK  $q \in x$  such that  $q \notin y$
  - $\langle 2 \rangle 4$ . Let:  $r \in y$ Prove:  $r \in x$
  - $\langle 2 \rangle 5. \ q \not\leq r$

PROOF: Since y is closed downwards.

- $\langle 2 \rangle 6$ . r < q
- $\langle 2 \rangle 7. \ r \in x$

PROOF: Since x is closed downwards.

**Proposition 15.3.8.** Any bounded nonempty subset of  $\mathbb{R}$  has a least upper bound.

#### Proof:

- $\langle 1 \rangle 1$ . Let: A be a bounded nonempty subset of  $\mathbb{R}$ .
- $\langle 1 \rangle 2$ .  $\bigcup A$  is a Dedekind cut.
  - $\langle 2 \rangle 1. \bigcup A \neq \emptyset$ 
    - $\langle 3 \rangle 1$ . Pick $x \in A$
    - $\langle 3 \rangle 2$ . Pick  $q \in x$
    - $\langle 3 \rangle 3. \ q \in \bigcup A$
  - $\langle 2 \rangle 2$ .  $\bigcup A \neq \mathbb{Q}$ 
    - $\langle 3 \rangle 1$ . Pick an upper bound u for A
    - $\langle 3 \rangle 2$ . Pick  $q \notin u$ Prove:  $q \notin \bigcup A$
    - $\langle 3 \rangle 3$ . Assume: for a contradiction  $q \in \bigcup A$
    - $\langle 3 \rangle 4$ . PICK  $x \in A$  such that  $q \in x$
    - $\langle 3 \rangle 5. \ x \leq u$
    - $\langle 3 \rangle 6. \ q \in u$
    - $\langle 3 \rangle$ 7. Q.E.D.

PROOF: This is a contradiction.

- $\langle 2 \rangle 3$ .  $\bigcup A$  is closed downwards.
  - $\langle 3 \rangle 1$ . Let:  $q \in \bigcup A$  and r < q
  - $\langle 3 \rangle 2$ . PICK  $x \in A$  such that  $q \in x$
  - $\langle 3 \rangle 3. \ r \in x$
  - $\langle 3 \rangle 4. \ r \in \bigcup A$
- $\langle 2 \rangle 4$ .  $\bigcup A$  has no greatest element.
  - $\langle 3 \rangle 1$ . Let:  $q \in \bigcup A$
  - $\langle 3 \rangle 2$ . PICK  $x \in A$  such that  $q \in x$
  - $\langle 3 \rangle 3$ . Pick  $r \in x$  such that q < r
  - $\langle 3 \rangle 4. \ r \in \bigcup A$
- $\langle 1 \rangle 3$ .  $\bigcup A$  is an upper bound for A.

PROOF: For all  $x \in A$  we have  $x \subseteq \bigcup A$ .

 $\langle 1 \rangle 4$ . For any upper bound u for  $\bigcup A$  we have  $\bigcup A \leq u$ .

PROOF: If  $\forall x \in A.x \subseteq u$  we have  $\bigcup A \subseteq u$ .

**Definition 15.3.9** (Addition). Define addition + on the reals by

$$x + y := \{q + r \mid q \in x, r \in y\}$$
.

We prove this is well-defined.

#### Proof:

 $\langle 1 \rangle 1$ . Let:  $x, y \in \mathbb{R}$ 

PROVE: X + y is a Dedekind cut.

```
\langle 1 \rangle 2. \ x + y \neq \emptyset
   PROOF: Pick q \in x and r \in y; then q + r \in x + y.
\langle 1 \rangle 3. \ x + y \neq \mathbb{Q}
   \langle 2 \rangle 1. PICK q \notin x and r \notin y
           PROVE: q + r \notin x + y
   \langle 2 \rangle 2. Assume: for a contradiction q + r \in x + y
   \langle 2 \rangle 3. Pick q' \in x and r' \in y such that q + r = q' + r'
   \langle 2 \rangle 4. q' < q and r' < r
   \langle 2 \rangle 5. q' + r' < q + r
   \langle 2 \rangle 6. Q.E.D.
      PROOF: This is a contradiction.
\langle 1 \rangle 4. x + y is closed downwards.
   \langle 2 \rangle 1. Let: q \in x and r \in y
   \langle 2 \rangle 2. Let: s < q + r
            PROVE: s \in x + y
   \langle 2 \rangle 3. \ s - r < q
   \langle 2 \rangle 4. \ s - r \in x
   \langle 2 \rangle 5. s = (s - r) + r \in x + y
\langle 1 \rangle 5. x + y has no greatest element.
   \langle 2 \rangle 1. Let: q \in x and r \in y
            PROVE: There exists s \in x + y such that q + r < s
   \langle 2 \rangle 2. Pick q' \in x and r' \in y such that q < q' and r < r'
   \langle 2 \rangle 3. \ \ q + r < q' + r' \in x + y
```

**Proposition 15.3.10.** Addition on the reals agrees with addition on the rationals.

```
PROOF: \begin{split} &\langle 1 \rangle 1. \text{ Let: } q,r \in \mathbb{Q} \\ &\langle 1 \rangle 2. \ q \downarrow + r \downarrow \subseteq (q+r) \downarrow \\ &\text{PROOF: If } s_1 < q \text{ and } s_2 < r \text{ then } s_1 + s_2 < q + r. \\ &\langle 1 \rangle 3. \ (q+r) \downarrow \subseteq q \downarrow + r \downarrow \\ &\langle 2 \rangle 1. \text{ Let: } s < q + r \\ &\langle 2 \rangle 2. \ s - r < q \\ &\langle 2 \rangle 3. \text{ PICK } t \text{ such that } s - r < t < q \\ &\langle 2 \rangle 4. \ s - t < r \\ &\langle 2 \rangle 5. \ s = t + (s-t) \in q \downarrow + r \downarrow \end{split}
```

Proposition 15.3.11. Addition is associative.

$$x + (y + z) = \{q + r \mid q \in x, r \in y + z\}$$

$$= \{q + s_1 + s_2 \mid q \in x, s_1 \in y, s_2 \in z\}$$

$$= \{r + s_2 \mid r \in x + y, s_2 \in z\}$$

$$= (x + y) + z$$

Proposition 15.3.12. Addition is commutative.

Proof:

$$x + y = \{q + r \mid q \in x, r \in y\}$$

$$= \{r + q \mid r \in y, q \in x\}$$

$$= y + x$$

**Proposition 15.3.13.** For any  $x \in \mathbb{R}$  we have x + 0 = x.

Proof:

 $\langle 1 \rangle 1$ .  $x + 0 \subseteq x$ 

PROOF: If  $q \in x$  and r < 0 then q + r < q so  $q + r \in x$ .

- $\langle 1 \rangle 2. \ x \subseteq x + 0$ 
  - $\langle 2 \rangle 1$ . Let:  $q \in x$
  - $\langle 2 \rangle 2$ . Pick  $r \in x$  such that q < r.

PROOF: x has no greatest element.

- $\langle 2 \rangle 3. \ q-r < 0$
- $(2)4. \ q = r + (q r) \in x + 0$

**Definition 15.3.14.** For  $x \in \mathbb{R}$ , define  $-x := \{q \in \mathbb{Q} \mid \exists r > q. -r \notin x\}$ .

**Proposition 15.3.15.** For all  $x \in \mathbb{R}$  we have  $-x \in \mathbb{R}$ .

Proof:

- $\langle 1 \rangle 1$ . Let:  $x \in \mathbb{R}$
- $\langle 1 \rangle 2$ .  $-x \neq \emptyset$ 
  - $\langle 2 \rangle 1$ . Pick  $s \notin x$
  - $\langle 2 \rangle 2$ .  $-s-1 \in -x$
- $\langle 1 \rangle 3. -x \neq \mathbb{Q}$ 
  - $\langle 2 \rangle 1$ . Pick  $s \in x$

Prove:  $-s \notin -x$ 

- $\langle 2 \rangle 2$ . Assume: for a contradiction  $-s \in -x$
- $\langle 2 \rangle 3$ . PICK r > -s such that  $-r \notin x$
- $\langle 2 \rangle 4$ . -r < s
- $\langle 2 \rangle 5$ . Q.E.D.

PROOF: This contradicts the fact that x is closed downwards.

 $\langle 1 \rangle 4$ . -x is closed downwards.

PROOF: Immediate from definition.

- $\langle 1 \rangle 5$ . -x has no greatest element.
  - $\langle 2 \rangle 1$ . Let:  $q \in -x$
  - $\langle 2 \rangle 2$ . Pick r > q such that  $-r \notin x$
  - $\langle 2 \rangle 3$ . Pick s such that q < s < r
- $\langle 2 \rangle 4. \ s \in -x$

**Lemma 15.3.16.** Let p be a positive rational number. For any real number x, there exists a rational  $q \in x$  such that  $p + q \notin x$ .

```
Proof:
\langle 1 \rangle 1. Pick q_0 \in x
\langle 1 \rangle 2. There exists k \in \mathbb{N} such that q_0 + kp \notin x
    \langle 2 \rangle 1. Pick q_1 \notin x
    \langle 2 \rangle 2. PICK k \in \mathbb{N} such that q_1 - q_0 < pk
       Proof: Proposition 15.2.11.
    \langle 2 \rangle 3. \ q_1 < q_0 + kp
    \langle 2 \rangle 4. \ q_0 + kp \notin x
\langle 1 \rangle 3. Let: k be the least natural number such that q_0 + kp \notin x
\langle 1 \rangle 4. \ k \neq 0
   Proof: \langle 1 \rangle 1
\langle 1 \rangle 5. Let: q = q_0 + (k-1)p
\langle 1 \rangle 6. \ q \in x \text{ and } q + p \notin x.
Proposition 15.3.17. For every real x we have x + (-x) = 0.
Proof:
\langle 1 \rangle 1. Let: x be a real number.
\langle 1 \rangle 2. x + (-x) \subseteq 0
    \langle 2 \rangle 1. Let: q_1 \in x and q_2 \in -x
    \langle 2 \rangle 2. PICK r > q_2 such that -r \notin x
    \langle 2 \rangle 3. \ q_1 < -r
    \langle 2 \rangle 4. r < -q_1
    \langle 2 \rangle 5. q_2 < -q_1
    \langle 2 \rangle 6. \ q_1 + q_2 < 0
\langle 1 \rangle 3. \ 0 \subseteq x + (-x)
    \langle 2 \rangle 1. Let: p < 0
    \langle 2 \rangle 2. 0 < -p
    \langle 2 \rangle 3. Pick q \in x such that q - p/2 \notin x
       Proof: Lemma 15.3.16.
    \langle 2 \rangle 4. Let: s = p/2 - q
    \langle 2 \rangle 5. -s \notin x
    \langle 2 \rangle 6. \ p - q < s
    \langle 2 \rangle 7. \ p-q \in -x
    \langle 2 \rangle 8. \ p \in x + (-x)
```

Corollary 15.3.17.1. The reals form an Abelian group under addition.

**Proposition 15.3.18.** For any reals x, y and z, we have x < y if and only if x + z < y + z.

```
\begin{split} &\langle 1 \rangle 1. \  \, \forall x,y,z \in \mathbb{R}. x \leq y \Rightarrow x+z \leq y+z \\ &\langle 2 \rangle 1. \  \, \text{Let:} \  \, x,y,z \in \mathbb{R} \\ &\langle 2 \rangle 2. \  \, \text{Assume:} \  \, x \leq y \\ &\langle 2 \rangle 3. \  \, \text{For all} \, \, q \in x \, \, \text{and} \, \, r \in z \, \, \text{we have} \, \, q+r \in y+z \end{split}
```

 $\langle 1 \rangle 2$ .  $\forall x, y, z \in \mathbb{R}.x + z = y + z \Leftrightarrow x = y$ PROOF: Proposition 12.1.4.

 $\langle 1 \rangle 3. \ \forall x, y, z \in \mathbb{R}. x < y \Rightarrow x + z < y + z$ 

 $\langle 1 \rangle 4$ . Q.E.D.

Proof: Proposition 5.2.6.

Γ

**Definition 15.3.19** (Absolute Value). The *absolute value* of a real number x is defined to be

$$|x| = \begin{cases} x & \text{if } 0 \le x \\ -x & \text{if } x < 0 \end{cases}$$

**Definition 15.3.20** (Multiplication). Define *multiplication*  $\cdot$  on  $\mathbb{R}$  as follows:

• If x and y are non-negative then

$$xy = 0 \cup \{rs \mid 0 \le r \in x \land 0 \le s \in y\} .$$

 $\bullet$  If x and y are both negative then

$$xy = (-x)(-y) .$$

• If one of x and y is negative and one is non-negative then

$$xy = -(|x||y|) .$$

We prove this is well-defined.

Proof:

 $\langle 1 \rangle 1$ . Let: x and y be non-negative reals.

PROVE: xy is real.

 $\langle 1 \rangle 2$ .  $xy \neq \emptyset$ 

PROOF: Since  $-1 \in xy$ .

 $\langle 1 \rangle 3. \ xy \neq \mathbb{Q}$ 

 $\langle 2 \rangle 1$ . Pick  $r \notin x$  and  $s \notin y$ 

Prove:  $rs \notin xy$ 

 $\langle 2 \rangle 2$ .  $0 \le r$  and  $0 \le s$ 

PROOF: Since  $0 \subseteq x$  and  $0 \subseteq y$ .

- $\langle 2 \rangle 3$ . Assume: for a contradiction  $rs \in xy$
- $\langle 2 \rangle 4$ . Pick r' and s' such that  $0 \leq r' \in x$ ,  $0 \leq s' \in y$  and rs = r's'
- $\langle 2 \rangle 5. \ r' < r$
- $\langle 2 \rangle 6. \ s' < s$
- $\langle 2 \rangle 7$ . r's' < rs
- $\langle 2 \rangle 8$ . Q.E.D.

PROOF: This is a contradiction.

- $\langle 1 \rangle 4$ . xy is closed downwards.
  - $\langle 2 \rangle 1$ . Let:  $q \in xy$  and r < q

```
\langle 2 \rangle 2. Case: q \in 0
      PROOF: Then r < q < 0 so r \in xy
   \langle 2 \rangle 3. Case: q = s_1 s_2 where 0 \le s_1 \in x and 0 \le s_2 \in y
      \langle 3 \rangle 1. Assume: w.l.o.g. 0 \le r
      \langle 3 \rangle 2. 0 < s_1 and 0 < s_2
      \langle 3 \rangle 3. \ r/s_2 < s_1
      \langle 3 \rangle 4. \ r/s_2 \in x
      \langle 3 \rangle 5. r = (r/s_2)s_2 \in xy
\langle 1 \rangle 5. xy has no greatest element.
   \langle 2 \rangle 1. Let: q \in xy
   \langle 2 \rangle 2. Case: q \in 0
      Proof: q < q/2 \in 0
   \langle 2 \rangle 3. Case: q = rs where 0 \le r \in x and 0 \le s \in y
      \langle 3 \rangle 1. Pick r' and s' with r < r' \in x and s < s' \in y
      \langle 3 \rangle 2. q < r's' \in xy
П
Proposition 15.3.21. Multiplication is commutative.
Proof: Immediate from definition.
Proposition 15.3.22. Multiplication is associative.
Proof:
\langle 1 \rangle 1. For non-negative reals x, y and z, we have x(yz) = (xy)z
   PROOF: It computes to 0 \cup \{qrs \mid 0 \le q \in x, 0 \le r \in y, 0 \le s \in z\}.
\langle 1 \rangle 2. For all reals x, y and z, we have x(yz) = (xy)z
   PROOF: It is equal to |x||y||z| if an even number of them are negative, and
   -(|x||y||z|) otherwise.
Proposition 15.3.23. Multiplication distributes over addition.
\langle 1 \rangle 1. For all non-negative reals x, y and z, we have x(y+z) = xy + xz
   \langle 2 \rangle 1. Let: x, y and z be non-negative reals.
   \langle 2 \rangle 2. x(y+z) \subseteq xy+xz
      \langle 3 \rangle 1. Let: q \in x(y+z)
      \langle 3 \rangle 2. Case: q < 0
         PROOF: Then we have q/2 \in xy and q/2 \in xz so q \in xy + xz.
      \langle 3 \rangle 3. Case: q = rs where 0 \le r \in x and 0 \le s \in y + z
         \langle 4 \rangle 1. PICK s_1 \in y and s_2 \in z such that s = s_1 + s_2
         \langle 4 \rangle 2. \ rs_1 \in xy
            PROOF: If s_1 < 0 then rs_1 < 0 so rs_1 \in xy. If 0 \le s_1 then we also
            have rs_1 \in xy.
         \langle 4 \rangle 3. \ rs_2 \in xz
            PROOF: Similar.
         \langle 4 \rangle 4. \ \ q \in xy + xz
```

PROOF: Since  $q = rs_1 + rs_2$ .

- $\langle 2 \rangle 3. \ xy + xz \subseteq x(y+z)$ 
  - $\langle 3 \rangle 1$ . Let:  $q \in xy$  and  $r \in xz$ .

PROVE:  $q + r \in x(y + z)$ 

 $\langle 3 \rangle 2$ . Case: q < 0 and r < 0

PROOF: Then q + r < 0 so  $q + r \in x(y + z)$ .

- $\langle 3 \rangle 3$ . Case: q < 0 and  $r = r_1 r_2$  where  $0 \le r_1 \in x$  and  $0 \le r_2 \in z$ 
  - $\langle 4 \rangle 1. \ q + r < r$
  - $\langle 4 \rangle 2. \ q + r \in xz$
  - $\langle 4 \rangle 3$ . Assume: w.l.o.g.  $0 \le q + r$

PROOF: Otherwise  $q + r \in x(y + z)$  immediately.

- $\langle 4 \rangle 4$ . PICK  $s_1, s_2$  with  $0 \leq s_1 \in x$ ,  $0 \leq s_2 \in y$  and  $q + r = s_1 s_2$
- $\langle 4 \rangle 5. \ s_2 \in y + z$

PROOF: Since  $0 \in z$  so  $s_2 = s_2 + 0 \in y + z$ .

- $\langle 4 \rangle 6. \ q+r \in x(y+z)$
- $\langle 3 \rangle 4$ . Case:  $q = q_1 q_2$  where  $0 \le q_1 \in x$  and  $0 \le q_2 \in y$  and r < 0 Proof: Similar.
- $\langle 3 \rangle 5.$  Case:  $q=q_1q_2$  where  $0\leq q_1\in x$  and  $0\leq q_2\in y$  and  $r=r_1r_2$  where  $0\leq r_1\in x$  and  $0\leq r_2\in z$ 
  - $\langle 4 \rangle 1$ . Assume: w.l.o.g.  $q_1 \leq r_1$
  - $\langle 4 \rangle 2. \ \ q + r \le r_1(q_2 + r_2) \in x(y + z)$
- $\langle 1 \rangle$ 2. For any negative real x and non-negative reals y and z, we have x(y+z)=xy+xz

Proof:

$$x(y+z) = -(-x)(y+z) = -((-x)y + (-x)z)$$
 ( $\langle 1 \rangle 1$ )  
=  $-((-x)y) - ((-x)z)$   
=  $xy + xz$ 

- $\langle 1 \rangle$ 3. For any non-negative real x and reals y and z with one negative and one non-negative, we have x(y+z)=xy+xz
  - $\langle 2 \rangle 1$ . Assume: w.l.o.g. y is negative and z is non-negative.
  - $\langle 2 \rangle 2$ . Case:  $0 \le y + z$

$$xy + xz = xy + x(-y + y + z)$$

$$= -(x(-y)) + x(-y + y + z)$$

$$= -(x(-y)) + x(-y) + x(y + z)$$

$$= x(y + z)$$
(\langle 1\rangle 1)

- $\langle 2 \rangle 3$ . Case: y + z < 0
  - $\langle 3 \rangle 1. -y z > 0$
  - $\langle 3 \rangle 2$ . -y = z y z
  - $\langle 3 \rangle 3$ . xy + xz = x(y+z)

Proof:

$$xy + xz = -(x(-y)) + xz$$

$$= -(x(z - y - z)) + xz$$

$$= -(xz + x(-y - z)) + xz \qquad (\langle 1 \rangle 1)$$

$$= -xy - x(-y - z) + xz$$

$$= -x(-y - z)$$

$$= x(y + z)$$

 $\langle 1 \rangle 4$ . For any non-negative real x and negative reals y and z, we have x(y+z)=xy+xz

Proof:

$$x(y+z) = -x(-y-z)$$

$$= -(x(-y) + x(-z))$$

$$= -x(-y) - x(-z)$$

$$= xy + xz$$

$$(\langle 1 \rangle 1)$$

- $\langle 1 \rangle$ 5. For any negative real x and reals y and z with one negative and one non-negative, we have x(y+z)=xy+xz
  - $\langle 2 \rangle 1$ . Assume: w.l.o.g. y is negative and z is non-negative.
  - $\langle 2 \rangle 2$ . Case:  $0 \le y + z$

Proof:

$$x(y+z) = -((-x)(y+z))$$

$$= -((-x)y + (-x)z)$$

$$= -((-x)y) - ((-x)z)$$

$$= (-x)(-y) - ((-x)z)$$

$$= xy + xz$$
(\langle 1\rangle 3)

 $\langle 2 \rangle 3$ . Case: y + z < 0

Proof:

$$x(y+z) = (-x)(-y-z) = (-x)(-y) + (-x)(-z) = xy + xz$$
 (\langle 1\rangle 3)

 $\langle 1 \rangle$ 6. For any negative reals x, y and z, we have x(y+z) = xy + xz Proof:

$$x(y+z) = (-x)(-y-z)$$

$$= (-x)(-y) + (-x)(-z)$$

$$= xy + xz$$

$$(\langle 1 \rangle 1)$$

**Proposition 15.3.24.** For any real x we have x1 = x.

- $\langle 1 \rangle 1$ . Case:  $0 \le x$ 
  - $\langle 2 \rangle 1. \ x1 \subseteq x$ 
    - $\langle 3 \rangle 1$ . Let:  $q \in x1$

⟨3⟩2. Case: 
$$q < 0$$
Proof: Then  $q \in x$  because  $0 \le x$ .
⟨3⟩3.  $q = rs$  where  $0 \le r \in x$  and  $0 \le s < 1$ 
Proof: Then  $q < r$  so  $q \in x$ .
⟨2⟩2.  $x \subseteq x1$ 
⟨3⟩1. Let:  $q \in x$ 
⟨3⟩2. Assume: w.l.o.g.  $0 \le q$ 
⟨3⟩3. Pick  $r$  such that  $q < r \in x$ 
⟨3⟩4.  $0 \le q/r < 1$ 
⟨3⟩5.  $q = r(q/r) \in x1$ 
⟨1⟩2. Case:  $x < 0$ 
Proof:
$$x1 = -((-x)1)$$

$$= x$$

**Lemma 15.3.25.** Let  $x \in \mathbb{R}$  and c be a positive rational. Then there exists  $a \in x$  and a non-least rational upper bound b for x such that b - a = c.

#### PROOF:

- (1)1. PICK  $a_1 \in x$  such that if x has a rational supremum s then  $a_1 > s c$
- $\langle 1 \rangle 2$ . There exists a natural number n such that  $a_1 + nc$  is an upper bound for x.
  - $\langle 2 \rangle 1$ . PICK a non-least upper bound  $b_1$  for x.
  - $\langle 2 \rangle 2$ . PICK a natural number n such that  $nc > b_1 a_1$

Proof: Proposition 15.2.11.

- $\langle 2 \rangle 3$ .  $a_1 + nc > b_1$
- $\langle 2 \rangle 4$ .  $a_1 + nc$  is an upper bound for x.
- $\langle 1 \rangle 3$ . Let: k be the least natural number such that  $a_1 + kc$  is an upper bound for x.
- $\langle 1 \rangle 4. \ a_1 + (k-1)c \in x$
- $\langle 1 \rangle 5$ .  $a_1 + kc$  is not the supremum of x.
  - $\langle 2 \rangle 1$ . Assume: for a contradiction  $a_1 + kc$  is the supremum of x.
  - $\langle 2 \rangle 2$ .  $a_1 > a_1 + (k-1)c$

Proof:  $\langle 1 \rangle 1$ 

 $\langle 2 \rangle 3$ . Q.E.D.

PROOF: This is a contradiction.

- $\langle 1 \rangle 6$ . Let:  $a = a_1 + (k-1)c$
- $\langle 1 \rangle 7$ . Let:  $b = a_1 + kc$
- $\langle 1 \rangle 8. \ b-a=c$

 $\Box$ 

**Proposition 15.3.26.** For any non-zero real x, there exists a real y such that xy = 1.

```
\langle 1 \rangle 1. Case: 0 < x
    \langle 2 \rangle 1. Let: y = \{q \in \mathbb{Q} \mid q \leq 0\} \cup \{u^{-1} \mid u \text{ is an upper bound for } x \text{ but not the supremum of } x\}
   \langle 2 \rangle 2. y is a real number.
       \langle 3 \rangle 1. \ y \neq \emptyset
           PROOF: Since 0 \in y.
       \langle 3 \rangle 2. \ y \neq \mathbb{Q}
           \langle 4 \rangle 1. PICK q \in x such that 0 < q
           \langle 4 \rangle 2. \ q^{-1} \notin y
       \langle 3 \rangle 3. y is closed downwards.
           \langle 4 \rangle 1. Let: q \in y and r < q
                    Prove: r \in y
           \langle 4 \rangle 2. Assume: w.l.o.g. 0 < r
           \langle 4 \rangle 3. q^{-1} is a non-least upper bound for x. \langle 4 \rangle 4. q^{-1} < r^{-1}
           \langle 4 \rangle 5. r^{-1} is a non-least upper bound for x.
           \langle 4 \rangle 6. \ r \in y
       \langle 3 \rangle 4. y has no greatest element.
           \langle 4 \rangle 1. Let: q \in y
                    PROVE: There exists r \in y such that q < r
           \langle 4 \rangle 2. Case: q \leq 0
               \langle 5 \rangle 1. PICK a non-least upper bound u for x.
               \langle 5 \rangle 2. \ q < u^{-1} \in x
           \langle 4 \rangle 3. Case: q = u^{-1} where u is a non-least upper bound for x.
               \langle 5 \rangle1. PICK a non-least upper bound v with v < u
               \langle 5 \rangle 2. \ u^{-1} < v^{-1} \in y
    \langle 2 \rangle 3. \ 0 < y
    \langle 2 \rangle 4. xy \subseteq 1
       \langle 3 \rangle 1. Let: q \in xy
       \langle 3 \rangle 2. Assume: w.l.o.g. 0 < q
       \langle 3 \rangle 3. Pick 0 < r \in x and 0 < s \in y such that q = rs
       \langle 3 \rangle 4. \ s^{-1} is a non-least upper bound for x
       \langle 3 \rangle 5. \ r < s^{-1}
       \langle 3 \rangle 6. \ rs < 1
    \langle 2 \rangle 5. 1 \subseteq xy
       \langle 3 \rangle 1. Let: q < 1
                Prove: q \in xy
       \langle 3 \rangle 2. Assume: w.l.o.g. 0 < q
       \langle 3 \rangle 3. PICK a_1 with 0 < a_1 \in x
       \langle 3 \rangle 4. \ (1-q)a_1 > 0
       \langle 3 \rangle 5. Pick a \in x and a non-least upper bound w of x such that w - a =
                (1-q)a_1
           PROOF: Lemma 15.3.25.
       \langle 3 \rangle 6. \ \ w - a < (1 - q)w
       \langle 3 \rangle 7. qw < a
       \langle 3 \rangle 8. \ w < a/q
       \langle 3 \rangle 9. a/q is a non-least upper bound for x
```

 $\langle 3 \rangle 10. \ q/a \in y$   $\langle 3 \rangle 11. \ q \in xy$   $\langle 1 \rangle 2. \ \text{Case:} \ x < 0$   $\langle 2 \rangle 1. \ \text{Pick} \ y \ \text{such that} \ (-x)y = 1$   $\text{Proof:} \ \langle 1 \rangle 1$   $\langle 2 \rangle 2. \ x(-y) = 1$ 

**Proposition 15.3.27.** For real numbers x, y and z, if 0 < z then x < y if and only if xz < yz.

#### Proof:

- $\langle 1 \rangle 1$ . For any real numbers x, y and z, if 0 < z and x < y then xz < yz
  - $\langle 2 \rangle 1$ . Let: x, y and z be real numbers.
  - $\langle 2 \rangle 2$ . Assume: 0 < z and x < y.
  - $\langle 2 \rangle 3. \ y = x + (y x)$
  - $\langle 2 \rangle 4$ . y x > 0
  - $\langle 2 \rangle 5$ . (y-x)z > 0
  - $\langle 2 \rangle 6. \ yz > xz$

Proof:

$$yz = (x + (y - x))z$$
$$= xz + (y - x)z$$
$$> xz$$

 $\langle 1 \rangle 2$ . For any real numbers x, y and z, if 0 < z and xz < yz then x < y PROOF: Proposition 5.2.6.

Corollary 15.3.27.1. The real numbers form a complete ordered field.

#### Proposition 15.3.28.

$$(0,1) \approx \mathbb{R}$$

PROOF: The function  $f(x) = (2x-1)/(x-x^2)$  is a bijection between (0,1) and  $\mathbb{R}$ .  $\square$ 

#### Proposition 15.3.29.

$$|\mathbb{R}| = 2^{\aleph_0}$$

#### Proof:

 $\langle 1 \rangle 1. \ (0,1) \leqslant 2^{\mathbb{N}}$ 

PROOF: The function H where H(x)(n) is the nth binary digit of the binary expansion of x is an injection.

 $\langle 1 \rangle 2. \ 2^{\mathbb{N}} \preccurlyeq \mathbb{R}$ 

PROOF: Map f to the real number in [0,1/9] whose n+1st decimal digit is f(n).

**Proposition 15.3.30.** The set of algebraic numbers is countable.

Proof:	There	are o	countably	many	integer	polynor	nials,	$\operatorname{each}$	with	finitely	many
roots.											

Corollary 15.3.30.1. There are uncountably many transcendental numbers.

**Proposition 15.3.31.** Let A be a set of disks in the plane, no two of which intersect. Then A is countable.

PROOF: Every circle includes a point with rational coordinates. Define  $f:\{q\in\mathbb{Q}^2\mid\exists C\in A.q\in C\}\rightarrow A$  by f(q)=C iff  $q\in C$ . Then f is surjective.  $\square$ 

**Proposition 15.3.32.** There exists an uncountable set of circles in the plane that do not intersect.

Proof: The set of all circles with origin O is uncountable.  $\square$ 

# Chapter 16

# Complex Analysis

**Definition 16.0.1.** For  $p \ge 1$ , let  $l^p$  be the set of all sequences of complex numbers  $(x_n)$  such that  $\sum_{n=1}^{\infty} |x_n|^p < \infty$ .

**Proposition 16.0.2.** If  $(x_n), (y_n) \in l^p$  then  $(x_n + y_n) \in l^p$ .

Proof:

$$\begin{array}{l} \langle 1 \rangle 1. \ \ \text{Let:} \ (x_n), (y_n) \in l^p \\ \langle 1 \rangle 2. \ \sum_{n=1}^{\infty} |x_n + y_n|^p \leq 2^p (\sum_{n=1}^{\infty} |x_n|^p + \sum_{n=1}^{\infty} |y_n|^p) \\ \text{PROOF:} \\ \langle 2 \rangle 1. \ \ \text{For all} \ n \in \mathbb{N} \ \text{we have} \ |x_n + y_n|^p \leq 2^p (|x_n|^p + |y_n|^p). \\ \text{PROOF:} \\ |x_n + y_n|^p \leq (|x_n| + |y_n|)^p \qquad \qquad \text{(Triangle Inequality)} \\ \leq (2 \max(|x_n|, |y_n|))^p \end{array}$$

 $\leq 2^p(|x_n|^p + |y_n|^p)$ 

**Theorem 16.0.3** (Hölder's Inequality). Let p and q be reals such that p > 1, q > 1 and 1/p + 1/q = 1. Let  $(x_n) \in l^p$  and  $(y_n) \in l^q$ . Then

$$\sum_{n} |x_n y_n| \le \left(\sum_{n} |x_n|^p\right)^{1/p} \left(\sum_{n} |y_n|^q\right)^{1/q}$$

Proof:

 $\langle 1 \rangle 1$ . Assume: w.l.o.g. neither  $(x_n)$  nor  $(y_n)$  are all zero.

 $\langle 1 \rangle 2$ . For  $0 \le x \le 1$  we have

$$x^{1/p} \le \frac{1}{p}x + \frac{1}{q} .$$

$$\langle 2 \rangle 2$$
,  $f'(x) = 1/p(1-x^{(1-p)/p})$ 

$$\langle 2 \rangle 3$$
.  $f'(x) > 0$  for all  $x \in [0, 1]$ 

 $x^{1/p} \leq \frac{1}{p}x + \frac{1}{q} .$   $\langle 2 \rangle 1.$  Let:  $f(x) = x/p + 1/q - x^{1/p}$   $\langle 2 \rangle 2.$   $f'(x) = 1/p(1 - x^{(1-p)/p})$   $\langle 2 \rangle 3.$   $f'(x) \geq 0$  for all  $x \in \mathbb{R}^n$   $\langle 2 \rangle 4.$   $f : \mathbb{R}^n$  $\langle 2 \rangle 4$ . f is a monotonically decreasing function on [0,1]

$$\langle 2 \rangle 5. \ f(0) = 1/q$$

$$\langle 2 \rangle 6. \ f(1) = 0$$

$$\langle 2 \rangle 7$$
.  $f(x) \geq 0$  for all  $x \in [0,1]$ 

 $\langle 1 \rangle 3$ . For any  $a, b \geq 0$  we have

$$ab \leq \frac{a^p}{p} + \frac{b^q}{q}$$
.

$$\langle 2 \rangle 1$$
. Case:  $a^p < b^q$ 

$$\langle 3 \rangle 1. \ ab^{-q/p} \le \frac{1}{p} \frac{a^p}{b^q} + \frac{1}{q}$$

$$\langle 3 \rangle 2$$
.  $ab^{1-q} \le \frac{1}{p} \frac{a^p}{b^q} + \frac{1}{q}$ 

 $\langle 2 \rangle 1. \text{ Case: } a^p \leq b^q$   $\langle 3 \rangle 1. \ ab^{-q/p} \leq \frac{1}{p} \frac{a^p}{b^q} + \frac{1}{q}$ Proof: Substituting  $x = a^p/b^q$  in  $\langle 1 \rangle 2$ .  $\langle 3 \rangle 2. \ ab^{1-q} \leq \frac{1}{p} \frac{a^p}{b^q} + \frac{1}{q}$ Proof: From  $\langle 3 \rangle 1$  since 1 - q = -q/p.  $\langle 3 \rangle 3. \ ab \leq \frac{a^p}{p} + \frac{b^q}{q}$ Proof: Multiplying  $\langle 3 \rangle 2$  by  $b^q$ .

$$\langle 3 \rangle 3$$
.  $ab \leq \frac{a^p}{p} + \frac{b^{\prime q}}{q}$ 

PROOF: Multiplying  $\langle 3 \rangle 2$  by  $b^q$ .

 $\langle 2 \rangle 2$ . Case:  $b^q \leq a^p$ 

Proof: Similar.

TROOF. Similar. 
$$\langle 1 \rangle 4$$
. For any integers  $1 \le j \le n$ , we have 
$$\frac{|x_j|}{\left(\sum_{k=1}^n |x_k|^p\right)^{1/p}} \frac{|y_j|}{\left(\sum_{k=1}^n |y_k|^q\right)^{1/q}} \le \frac{1}{p} \frac{|x_j|^p}{\sum_{k=1}^n |x_k|^p} + \frac{1}{q} \frac{|y_j|^q}{\sum_{k=1}^n |y_k|^q}$$
PROOF: From  $\langle 1 \rangle 3$  substituting 
$$a = \frac{|x_j|}{\left(\sum_{k=1}^n |x_k|^p\right)^{1/p}} \text{ and } b = \frac{|y_j|}{\left(\sum_{k=1}^n |y_k|^q\right)^{1/q}}$$
/1\(\frac{5}{5}\). For any positive integer  $n$  we have

$$a = \frac{|x_j|}{\left(\sum_{k=1}^n |x_k|^p\right)^{1/p}}$$
 and  $b = \frac{|y_j|}{\left(\sum_{k=1}^n |y_k|^q\right)^{1/q}}$ 

(1)5. For any positive integer 
$$n$$
 we have
$$\frac{\sum_{k=1}^{n} |x_k| |y_k|}{\left(\sum_{k=1}^{n} |x_k|^p\right)^{1/p} \left(\sum_{k=1}^{n} |y_k|^q\right)^{1/q}} \le 1$$
Proof:

Proof:

FROOF: 
$$\frac{\sum_{j=1}^{n} |x_j| |y_j|}{\left(\sum_{k=1}^{n} |x_k|^p\right)^{1/p} \left(\sum_{k=1}^{n} |y_k|^q\right)^{1/q}} \le \frac{1}{p} + \frac{1}{q} \quad \text{(Summing } \langle 1 \rangle 4 \text{ from } j = 1 \text{ to } n\text{)}$$

$$= 1$$

 $\langle 1 \rangle 6$ .

$$\sum_{n} |x_n y_n| \le \left(\sum_{n} |x_n|^p\right)^{1/p} \left(\sum_{n} |y_n|^q\right)^{1/q}$$

PROOF: Taking the limit  $n \to \infty$  in  $\langle 1 \rangle 5$ 

**Theorem 16.0.4** (Minkowski's Inequality). Let  $p \geq 1$ . Let  $(x_n), (y_n) \in l^p$ . Then

$$\left(\sum_{n=1}^{\infty} |x_n + y_n|^p\right)^{1/p} \le \left(\sum_{n=1}^{\infty} |x_n|^p\right)^{1/p} + \left(\sum_{n=1}^{\infty} |y_n|^p\right)^{1/p}$$

Proof:

 $\langle 1 \rangle 1$ . Case: p = 1

PROOF: This is just the Triangle Inequality.

 $\langle 1 \rangle 2$ . Case: p > 1

$$\langle 2 \rangle 1$$
. Let:  $q = p/(p-1)$ 

$$\langle 2 \rangle 2$$
.

$$\sum_{n=1}^{\infty} |x_n + y_n|^p \le \left(\sum_{n=1}^{\infty} |x_n|^p\right)^{1/p} \left(\sum_{n=1}^{\infty} |x_n + y_n|^{q(p-1)}\right)^{1/q} + \left(\sum_{n=1}^{\infty} |y_n|^p\right)^{1/p} \left(\sum_{n=1}^{\infty} |x_n + y_n|^{q(p-1)}\right)^{1/q}$$

Proof:

$$\langle 3 \rangle 1. \ (|x_n + y_n|^{p-1}) \in l^q$$
PROOF:

$$\sum_{n=1}^{\infty} |x_n + y_n|^{(p-1)q} = \sum_{n=1}^{\infty} |x_n + y_n|^p$$

$$< \infty \qquad (\text{Proposition 16.0.2})$$

 $\langle 3 \rangle 2$ . Q.E.D.

PROOF:
$$\sum_{n=1}^{\infty} |x_n + y_n|^p = \sum_{n=1}^{\infty} |x_n + y_n| |x_n + y_n|^{p-1}$$

$$\leq \sum_{n=1}^{\infty} |x_n| |x_n + y_n|^{p-1} + \sum_{n=1}^{\infty} |y_n| |x_n + y_n|^{p-1}$$

$$\leq \left(\sum_{n=1}^{\infty} |x_n|^p\right)^{1/p} \left(\sum_{n=1}^{\infty} |x_n + y_n|^{q(p-1)}\right)^{1/q}$$

$$+ \left(\sum_{n=1}^{\infty} |y_n|^p\right)^{1/p} \left(\sum_{n=1}^{\infty} |x_n + y_n|^{q(p-1)}\right)^{1/q}$$
(Hölder's Inequality,  $\langle 2 \rangle 2$ )

 $\langle 2 \rangle 3$ .

$$\sum_{n=1}^{\infty} |x_n + y_n|^p \le \left\{ \left( \sum_{n=1}^{\infty} |x_n|^p \right)^{1/p} + \left( \sum_{n=1}^{\infty} |y_n|^p \right)^{1/p} \right\} \left( \sum_{n=1}^{\infty} |x_n + y_n|^p \right)^{1/q}$$

 $\langle 3 \rangle 1. \ q(p-1) = p$ 

Proof:  $\langle 2 \rangle 2$ 

 $\langle 3 \rangle 2$ . Q.E.D.

PROOF: From  $\langle 2 \rangle 2$ ,  $\langle 3 \rangle 1$ .

# Part I Linear Algebra

# Chapter 17

# Vector Spaces

## 17.1 Vector Spaces

**Definition 17.1.1** (Vector Space). Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . A vector space over K is a triple  $(V, +, \cdot)$  such that:

- $\bullet$  V is a nonempty set, whose elemnts are called *vectors*;
- $\bullet \ +: V^2 \to V$
- $\bullet : K \times V \to V$

such that the following hold for all  $u, v, w \in V$  and  $\alpha, \beta \in K$ :

- 1. u + v = v + u
- 2. u + (v + w) = (u + v) + w
- 3. For every  $u, v \in V$  there exists  $w \in V$  such that u + w = v
- 4.  $\alpha(\beta v) = (\alpha \beta)v$
- 5.  $(\alpha + \beta)v = \alpha v + \beta v$
- 6.  $\alpha(u+v) = \alpha u + \alpha v$
- 7. 1v = v

Elements of K are called *scalars*.

We write real vector space for 'vector space over  $\mathbb{R}$ ', and complex vector space for 'vector space over  $\mathbb{C}$ '.

**Proposition 17.1.2.** Let K be either  $\mathbb{R}$  and  $\mathbb{C}$ . The set  $\{0\}$  is a vector space over K under the unique functions  $+: \{0\}^2 \to \{0\}, : K \times \{0\} \to \{0\}$ .

PROOF: Each axiom holds trivially because x = y holds for all  $x, y \in \{0\}$ .  $\square$ 

**Proposition 17.1.3.** The set  $\mathbb{R}$  is a real vector space under real addition and real multiplication.

PROOF: TODO — after we have proved these facts about  $\mathbb{R}$ .  $\square$ 

**Proposition 17.1.4.** The set  $\mathbb{C}$  is a real vector space under complex addition and complex multiplication.

PROOF: TODO

**Proposition 17.1.5.** The set  $\mathbb{C}$  is a complex vector space under complex addition and complex multiplication.

PROOF: TODO

**Proposition 17.1.6.** Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . Let  $\{V_i\}_{i\in I}$  be a family of vector spaces over K. Then  $\prod_{i\in I} V_i$  is a vector space over K under the operations given by

$${x_i}_{i \in I} + {y_i}_{i \in I} = {x_i + y_i}_{i \in I}$$
  
 $\alpha {x_i}_{i \in I} = {\alpha x_i}_{i \in I}$ 

PROOF: Each axiom follows from the corresponding axiom in  $V_i$ .

**Corollary 17.1.6.1.** Let V be a vector space over K. For any set I, we have  $V^I$  is a vector space over K.

**Corollary 17.1.6.2.** Let  $n \in \mathbb{Z}_+$ . Then  $\mathbb{R}^n$  is a real vector space, and  $\mathbb{C}^n$  is both a real and a complex vector space, under

$$(x_1, \dots, x_n) + (y_1, \dots, y_n) = (x_1 + y_1, \dots, x_n + y_n)$$
  
 $\lambda(x_1, \dots, x_n) = (\lambda x_1, \dots, \lambda x_n)$ 

**Proposition 17.1.7.** Let V be a vector space over K. Then there exists a unique  $0 \in V$  such that, for all  $v \in V$ , we have v + 0 = v.

#### PROOF:

- $\langle 1 \rangle 1$ . There exists  $0 \in V$  such that  $\forall v \in V.v + 0 = v$ 
  - $\langle 2 \rangle 1$ . Pick  $v \in V$
  - $\langle 2 \rangle 2$ . Pick  $0 \in V$  such that v + 0 = v

Proof: Axiom 3.

- $\langle 2 \rangle 3$ . For all  $u \in V$ , we have u + 0 = u
  - $\langle 3 \rangle 1$ . Let:  $u \in V$
  - $\langle 3 \rangle 2$ . Pick  $u' \in V$  such that v + u' = u

Proof: Axiom 3.

 $\langle 3 \rangle 3. \ u + 0 = u$ 

$$u + 0 = v + u' + 0 \tag{\langle 3 \rangle 2}$$

$$= v + u' \tag{2}2$$

$$=u$$
  $(\langle 3 \rangle 2)$ 

$$\langle 1 \rangle 2$$
. If  $0, 0' \in V$  are such that  $\forall v \in V.v + 0 = v$  and  $\forall v \in V.v + 0' = v$ , then  $0 = 0'$ .

- $\langle 2 \rangle 1$ . Let:  $0, 0' \in V$
- $\langle 2 \rangle 2$ . Assume:  $\forall v \in V.v + 0 = v$
- $\langle 2 \rangle 3$ . Assume:  $\forall v \in V.v + 0' = v$
- $\langle 2 \rangle 4. \ 0 = 0'$

$$0 = 0 + 0' \tag{\langle 2 \rangle 2}$$

$$=0' \qquad (\langle 2 \rangle 3)$$

**Proposition 17.1.8.** Let V be a vector space. For any  $v \in V$ , there exists a unique  $-v \in V$  such that v + (-v) = 0.

PROOF:

- $\langle 1 \rangle 1$ . Let:  $v \in V$
- $\langle 1 \rangle 2$ . There exists  $-v \in V$  such that v + (-v) = u

PROOF: Axiom 3.

- $\langle 1 \rangle 3$ . If v + x = 0 and v + y = 0 then x = y
  - $\langle 2 \rangle 1$ . Assume: v + x = 0
  - $\langle 2 \rangle 2$ . Assume: v + y = 0
  - $\langle 2 \rangle 3. \ x = y$

Proof:

$$x = x + 0$$
 (Proposition 17.1.7)  
 $= x + v + y$  ( $\langle 2 \rangle 2$ )  
 $= 0 + y$  ( $\langle 2 \rangle 1$ )  
 $= y$  (Proposition 17.1.7)

**Proposition 17.1.9.** Let V be a vector space. For any  $u, v \in V$ , there exists a unique  $u - v \in V$  such that v + (u - v) = u, namely u - v = u + (-v).

Proof:

- $\langle 1 \rangle 1$ . Let:  $u, v \in V$
- $\langle 1 \rangle 2. \ v + (u + (-v)) = u$

Proof:

$$v + u + (-v) = u + 0$$
 (Proposition 17.1.8)  
=  $u$  (Proposition 17.1.7)

 $\langle 1 \rangle 3$ . For all  $x \in V$ , if v + x = u then x = u + (-v).

- $\langle 2 \rangle 1$ . Let:  $x \in V$
- $\langle 2 \rangle 2$ . Assume: v + x = u
- $\langle 2 \rangle 3$ . x = u + (-v)

$$u + (-v) = v + x + (-v)$$
 ( $\langle 2 \rangle 2$ )  
=  $x + 0$  (Proposition 17.1.8)  
=  $x$  (Proposition 17.1.7)

П

**Proposition 17.1.10.** Let V be a vector space over K. Let  $u, v, w \in V$ . If u + v = u + w then v = w.

Proof:

$$\langle 1 \rangle 1$$
. Assume:  $u + v = u + w$ 

 $\langle 1 \rangle 2. \ v = w$ 

Proof:

$$v = v + 0$$
 (Proposition 17.1.7)  
 $= v + u + (-u)$  (Proposition 17.1.8)  
 $= w + u + (-u)$  ( $\langle 1 \rangle 1$ )  
 $= w + 0$  (Proposition 17.1.8)  
 $= w$  (Proposition 17.1.7)

**Proposition 17.1.11.** Let V be a vector space over K. Let  $\lambda \in K$ . Then  $\lambda 0 = 0$ .

Proof:

$$\langle 1 \rangle 1$$
.  $\lambda 0 + \lambda 0 = \lambda 0 + 0$ 

Proof:

$$\lambda 0 + \lambda 0 = \lambda (0 + 0)$$
 (Axiom 6)  
=  $\lambda 0$  (Proposition 17.1.7)

 $\langle 1 \rangle 2$ .  $\lambda 0 = 0$ 

Proof: Proposition 17.1.10.

**Proposition 17.1.12.** Let V be a vector space over K. Let  $\lambda \in K$  and  $v \in V$ . If  $\lambda v = 0$  then  $\lambda = 0$  or v = 0.

Proof:

- $\langle 1 \rangle 1$ . Assume:  $\lambda \neq 0$
- $\langle 1 \rangle 2$ . Assume:  $\lambda v = 0$
- $\langle 1 \rangle 3. \ v = 0$

Proof:

$$v = 1v$$
 (Axiom 7)  
 $= \lambda^{-1} \lambda v$   
 $= \lambda^{-1} 0$  ( $\langle 1 \rangle 2$ )  
 $= 0$ 

**Proposition 17.1.13.** Let V be a vector space over K. For all  $v \in V$  we have 0v = 0.

$$\langle 1 \rangle 1$$
.  $0v + 0 = 0v + 0v$ 

$$0v + 0 = 0v \qquad (Proposition 17.1.7)$$

$$= (0+0)v$$

$$= 0v + 0v \qquad (Axiom 5)$$

$$\langle 1 \rangle 2. \ 0v = 0$$

$$PROOF: Proposition 17.1.10, \langle 1 \rangle 1.$$

$$\square$$

$$Proposition 17.1.14. Let V be a vector space over K. Let v$$

**Proposition 17.1.14.** Let V be a vector space over K. Let  $v \in V$ . Then (-1)v = -v.

PROOF:  $\langle 1 \rangle 1. \ v + (-1)v = 0 \\ \text{PROOF:} \\ v + (-1)v = 1v + (-1)v \qquad \qquad \text{(Axiom 7)} \\ = (1 + (-1))v \qquad \qquad \text{(Axiom 5)} \\ = 0v \qquad \qquad = 0 \\ \langle 1 \rangle 2. \ \text{Q.E.D.} \\ \text{PROOF: Proposition 17.1.8.} \\ \sqcap$ 

## 17.2 Subspaces

**Definition 17.2.1** (Subspace). Let V be a vector space over K and  $U \subseteq V$ . Then U is a *subspace* of V iff  $\forall \alpha, \beta \in K. \forall u, v \in U. \alpha u + \beta v \in U$ . It is a *proper* subspace iff in addition  $U \neq V$ .

**Proposition 17.2.2.** Let V be a vector space over K and U a subspace of V. Then U is a vector space over K under the restrictions of the operations of V.

PROOF: Each of the axioms follows from the corresponding axiom in V. For axiom 3, we have if  $u, v \in U$  then  $v - u = 1v + (-1)u \in U$ .  $\square$ 

Proposition 17.2.3. Every vector space is a subspace of itself.

Proof: Trivial.

**Proposition 17.2.4.** Let  $\Omega$  be a subset of  $\mathbb{R}^N$ . Let  $\mathcal{C}(\Omega)$  be the set of all continuous functions  $\Omega \to \mathbb{C}$ . Then  $\mathcal{C}(\Omega)$  is a subspace of  $\mathbb{C}^{\Omega}$ .

PROOF: If  $f, g: \Omega \to \mathbb{C}$  are continuous then so is  $\alpha f + \beta g$ .  $\square$ 

**Proposition 17.2.5.** Let  $\Omega$  be an open set in  $\mathbb{R}^N$ . Let  $\mathcal{C}^k(\Omega)$  be the set of all continuous functions  $\Omega \to \mathbb{C}$  with continuous partial derivatives of order k. Then  $\mathcal{C}^k(\Omega)$  is a subspace of  $\mathbb{C}^{\Omega}$ .

PROOF: If  $f,g:\Omega\to\mathbb{C}$  have continuous partial derivatives of order k then so does  $\alpha f+\beta g$ .  $\square$ 

**Proposition 17.2.6.** Let  $\Omega$  be an open set in  $\mathbb{R}^N$ . Let  $\mathcal{C}^{\infty}(\Omega)$  be the set of all infinitely differentiable functions  $\Omega \to \mathbb{C}$ . Then  $\mathcal{C}^{\infty}(\Omega)$  is a subspace of  $\mathbb{C}^{\Omega}$ .

PROOF: If  $f, g: \Omega \to \mathbb{C}$  are infinitely differentiable then so is  $\alpha f + \beta g$ .  $\square$ 

**Proposition 17.2.7.** Let  $\Omega$  be an open set in  $\mathbb{R}^N$ . Let  $\mathcal{P}(\Omega)$  be the set of all polynomials in N variables considered as functions  $\Omega \to \mathbb{C}$ . Then  $\mathcal{P}(\Omega)$  is a subspace of  $\mathbb{C}^{\Omega}$ .

PROOF: If  $f, g: \Omega \to \mathbb{C}$  are polynomials in N variables then so is  $\alpha f + \beta g$ .  $\square$ 

**Proposition 17.2.8.** Let V be a vector space and  $U_1$ ,  $U_2$  subspaces of V. If  $U_1 \subseteq U_2$  then  $U_1$  is a subspace of  $U_2$ .

PROOF: Trivial.  $\square$ 

**Proposition 17.2.9.** Let V be a vector space over K. The intersection of a set of subspaces of V is a subspace of V.

#### Proof:

```
\begin{split} &\langle 1 \rangle 1. \text{ Let: } \mathcal{U} \text{ be a set of subspaces of } V. \\ &\langle 1 \rangle 2. \text{ Let: } u,v \in \bigcap \mathcal{U} \text{ and } \lambda,\mu \in K \\ &\langle 1 \rangle 3. \ \lambda u + \mu v \in \bigcap \mathcal{U} \\ &\langle 2 \rangle 1. \text{ Let: } \mathcal{U} \in \mathcal{U} \\ &\langle 2 \rangle 2. \ u,v \in \mathcal{U} \\ &\text{Proof: } \langle 1 \rangle 2,\ \langle 2 \rangle 1. \\ &\langle 2 \rangle 3. \ \lambda u + \beta v \in \mathcal{U} \\ &\text{Proof: } \langle 1 \rangle 1,\ \langle 1 \rangle 2,\ \langle 2 \rangle 1,\ \langle 2 \rangle 2. \\ &\Box \end{split}
```

**Proposition 17.2.10.** The set of all bounded complex sequences is a proper subspace of  $\mathbb{C}^{\mathbb{N}}$ .

PROOF: If  $(x_n)$  and  $(y_n)$  are bounded then so is  $(\lambda x_n + \mu y_n)$ .  $\square$ 

**Proposition 17.2.11.** The set of all convergent complex sequences is a proper subspace of the space of all bounded complex sequences.

PROOF: If  $(x_n)$  and  $(y_n)$  converge then so does  $(\lambda x_n + \mu y_n)$ .  $\square$ 

**Proposition 17.2.12.** The set  $l^p$  of all sequences  $(x_n)$  in  $\mathbb{C}$  such that  $\sum_n |x_n|^p < \infty$  is a subspace of  $\mathbb{C}^{\mathbb{N}}$ .

PROOF: It is closed under addition by Proposition 16.0.2, and it is easy to see that it is closed under scalar multiplication.  $\Box$ 

### 17.3 Linear Independence and Bases

**Definition 17.3.1** (Linear Combination). Let V be a vector space over K. Let  $v, v_1, \ldots, v_n \in V$ . Then v is a *linear combination* of  $v_1, \ldots, v_n$  iff there exist scalars  $\lambda_1, \ldots, \lambda_n \in K$  such that

$$v = \lambda_1 v_1 + \dots + \lambda_n v_n$$
.

**Definition 17.3.2** (Linearly Independent). Let V be a vector space over K. Let  $A \subseteq V$ . Then A is linearly independent iff, for all  $\lambda_1, \ldots, \lambda_n \in K$  and  $v_1, \ldots, v_n \in A$ , if  $\lambda_1 v_1 + \cdots + \lambda_n v_n = 0$  then  $\lambda_1 = \cdots = \lambda_n = 0$ .

**Definition 17.3.3** (Span). Let V be a vector space over K and  $A \subseteq V$ . The span of A, or the subspace of V spanned by A, is the set of all linear combinations of vectors in A.

**Proposition 17.3.4.** Let V be a vector space over K and  $A \subseteq V$ . Then span A is a subspace of V.

PROOF: Given  $\alpha, \beta \in K$  and  $\lambda_1 u_1 + \cdots + \lambda_m u_m, \mu_1 v_1 + \cdots + \mu_n v_n \in \operatorname{span} A$ , we have

$$\alpha(\lambda_1 u_1 + \dots + \lambda_m u_m) + \beta(\mu_1 v_1 + \dots + \mu_n v_n)$$

$$= \alpha \lambda_1 u_1 + \dots + \alpha \lambda_m u_m + \beta \mu_1 v_1 + \dots + \beta \mu_n v_n$$

$$\in \operatorname{span} A$$

**Definition 17.3.5** (Basis). Let V be a vector space over K and  $B \subseteq V$ . Then B is a *basis* for V iff B is linearly independent and span B = V.

**Definition 17.3.6** (Finite Dimensional). A vector space is *finite dimensional* iff there exists a finite basis; otherwise it is *infinite dimensional*.

**Proposition 17.3.7.** In a finite dimensional space, any two bases have the same size.

TODO

**Definition 17.3.8** (Dimension). The *dimension* of a finite dimensional vector space V, dim V, is the number of vectors in any basis.

**Proposition 17.3.9.** Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . Then  $K^n$  as a vector space over K has dimension n.

PROOF: The vectors with one component 1 and all other components 0 form a basis.  $\Box$ 

**Proposition 17.3.10.** As a real vector space,  $\mathbb{C}^n$  has dimension 2n.

PROOF: The vectors with one component either 1 or i and all other components 0 form a basis.  $\square$ 

**Proposition 17.3.11.** Let  $\Omega$  be a nonempty open set in  $\mathbb{R}^n$ . The space  $\mathcal{C}(\Omega)$  is infinite dimensional.

PROOF: Let  $\pi_1 : \mathbb{R}^n \to \mathbb{R}$  be the first projection. The functions  $1, \pi_1(x), \pi_1(x)^2, \pi_1(x)^3, \ldots$  form an infinite linearly independent set in  $\mathcal{C}(\Omega)$ .  $\square$ 

**Proposition 17.3.12.** The spaces  $C^k(\mathbb{R}^n)$  and  $C^{\infty}(\mathbb{R}^n)$  are infinite dimensional

Proof: The monomials 1, x,  $x^2$ , ... form an infinite linearly independent set.  $\sqcap$ 

## 17.4 Linear Mappings

**Definition 17.4.1** (Kernel). Let U and V be vector spaces and  $T: U \to V$ . The *kernel* of T is

$$\ker T := \{ u \in U \mid T(u) = 0 \}$$
.

**Definition 17.4.2** (Linear Mapping). Let U and V be vector spaces over K. A function  $L: U \to V$  is a linear mapping iff  $\forall x, y \in U. \forall \alpha, \beta \in K. L(\alpha x + \beta y) = \alpha L(x) + \beta L(y)$ .

**Proposition 17.4.3.** Let U and V be vector spaces over K. The set of linear mappings from U to V is a subspace of  $V^U$ .

## 17.5 Eigenvalues and Eigenvectors

**Definition 17.5.1** (Eigenvalue and Eigenvector). Let V be a vector space over K. Let  $A: V \to V$  be a linear transformation. Let  $v \in V$  and  $\lambda \in K$ . Then v is an eigenvector of A with eigenvalue  $\lambda$  iff  $A(v) = \lambda v$ .

## Chapter 18

# Normed Spaces

**Definition 18.0.1** (Norm). Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . Let V be a vector space over K. A *norm* on V is a function  $\| \ \| : V \to \mathbb{R}$  such that, for all  $u, v \in V$  and  $\lambda \in K$ :

- 1. If ||v|| = 0 then v = 0.
- $2. \|\lambda v\| = |\lambda| \|v\|$
- 3. (Triangle Inequality)  $||u+v|| \le ||u|| + ||v||$

A normed space over K is a pair (V, || ||) where V is a vector space over K and || || is a norm on V.

**Proposition 18.0.2.** In a normed space, ||0|| = 0.

PROOF: 
$$||0|| = |0|||0|| = 0$$
 by Axiom 2.  $\Box$ 

**Proposition 18.0.3.** Let V be a normed vector space over K. For all  $v \in V$  we have  $||v|| \ge 0$ .

Proof:

$$0 = ||0||$$
 (Proposition 18.0.2)  

$$= ||v - v||$$
 (Triangle Inequality)  

$$= 2||v||$$
 (Axiom 2)

**Proposition 18.0.4.** Let V be a normed space. Let  $u, v \in V$ . Then

$$|||u|| - ||v||| \le ||u - v||$$
.

Proof:

$$||u|| \le ||u - v|| + ||v||$$
 (Triangle Inequality)  

$$\therefore ||u|| - ||v|| \le ||u - v||$$
 (Triangle Inequality)  

$$= ||v - v|| + ||u||$$
 (Axiom 2)  

$$\therefore ||v|| - ||u|| \le ||u - v||$$

**Definition 18.0.5** (Euclidean Norm). The *Euclidean norm* on  $K^n$  is defined by

$$||(x_1,\ldots,x_n)|| = \sqrt{|x_1|^2 + \cdots + |x_n|^2}$$
.

**Proposition 18.0.6.** The Euclidean norm on  $K^n$  is a norm.

Proof:

$$\langle 1 \rangle 1$$
. If  $\|\vec{x}\| = 0$  then  $\vec{x} = \vec{0}$   
PROOF: If  $\sqrt{|x_1|^2 + \cdots + |x_n|^2} = 0$  then  $x_1 = \cdots = x_n = 0$ .  $\langle 1 \rangle 2$ .  $\|\lambda \vec{x}\| = |\lambda| \|\vec{x}\|$   
PROOF:

$$\|\lambda \vec{x}\| \sqrt{|\lambda x_1|^2 + \dots + |\lambda x_n|^2}$$

$$= \sqrt{|\lambda|^2 |x_1|^2 + \dots + |\lambda|^2 |x_n|^2}$$

$$= |\lambda| \sqrt{|x_1|^2 + \dots + |x_n|^2}$$

$$= |\lambda| \|\vec{x}\|$$

$$\langle 1 \rangle 3. \ \|\vec{u} + \vec{v}\| \le \|\vec{u}\| + \|\vec{v}\|$$

Proof:

$$\|\vec{u} + \vec{v}\|^{2} = |u_{1} + v_{1}|^{2} + \dots + |u_{n} + v_{n}|^{2}$$

$$= |u_{1}|^{2} + \dots + |u_{n}|^{2} + |v_{1}|^{2} + \dots + |v_{n}|^{2}$$

$$+ 2|u_{1}||v_{1}| + \dots + 2|u_{n}||v_{n}|$$

$$\leq \|\vec{u}\|^{2} + \|\vec{v}\|^{2} + 2|u_{1}v_{1} + \dots + u_{n}v_{n}|$$

$$\leq \|\vec{u}\|^{2} + \|\vec{v}\|^{2} + 2\|\vec{u}\|\|\vec{v}\| \qquad \text{(Cauchy-Schwarz)}$$

$$= (\|\vec{u}\| + \|\vec{v}\|)^{2}$$

Corollary 18.0.6.1. The absolute value function | | is a norm on K.

**Proposition 18.0.7.** The function  $\|\vec{x}\| = |x_1| + \cdots + |x_n|$  is a norm on  $\mathbb{C}^n$ .

$$\langle 1 \rangle 1$$
. If  $||\vec{x}|| = 0$  then  $\vec{x} = \vec{0}$   
PROOF: If  $|x_1| + \cdots + |x_n| = 0$  then  $x_1 = \cdots = x_n = 0$ .  $\langle 1 \rangle 2$ .  $||\lambda \vec{x}|| = |\lambda| ||\vec{x}||$ 

Proof:

$$\|\lambda \vec{x}\| |\lambda x_1| + \dots + |\lambda x_n|$$

$$= |\lambda| (|x_1| + \dots + |x_n|)$$

$$= |\lambda| \|\vec{x}\|$$

$$\langle 1 \rangle 3. \|\vec{u} + \vec{v}\| \le \|\vec{u}\| + \|\vec{v}\|$$
PROOF:
$$\|\vec{u} + \vec{v}\|^2 = |u_1 + v_1| + \dots + |u_n + v_n|$$

$$\le |u_1| + |v_1| + \dots + |u_n| + |v_n|$$

$$= \|\vec{u}\| + \|\vec{v}\|$$

**Proposition 18.0.8.** The function  $\|\vec{x}\| = \max(|x_1|, \dots, |x_n|)$  is a norm on  $\mathbb{C}^n$ .

Proof:

$$\begin{array}{l} \text{TROOF:} \\ \langle 1 \rangle 1. \text{ If } \|\vec{x}\| = 0 \text{ then } \vec{x} = \vec{0} \\ \text{PROOF:} \text{ If } \max(|x_1|, \dots, |x|n|) = 0 \text{ then } x_1 = \dots = x_n = 0. \\ \langle 1 \rangle 2. \ \|\lambda \vec{x}\| = |\lambda| \|\vec{x}\| \\ \text{PROOF:} \\ \|\lambda \vec{x}\| = \max(|\lambda x_1|, \dots, |\lambda x_n|) \\ &= |\lambda| \max(|x_1|, \dots, |x_n|) \\ &= |\lambda| \|\vec{x}\| \\ \langle 1 \rangle 3. \ \|\vec{u} + \vec{v}\| \leq \|\vec{u}\| + \|\vec{v}\| \\ \text{PROOF:} \\ \|\vec{u} + \vec{v}\| = \max(|u_1 + v_1|, \dots, |u_n + v_n|) \\ &\leq \max(|u_1| + |v_1|, \dots, |u_n| + |v_n|) \\ &\leq \max(|u_1|, \dots, |u_n|) + \max(|v_1|, \dots, |v_n|) \end{array}$$

**Definition 18.0.9** (Uniform Convergence Norm). Let  $\Omega$  be a closed bounded subset of  $\mathbb{R}^n$ . The *uniform convergence norm* on  $\mathcal{C}(\Omega)$  is the function defined by  $||f|| = \max_{x \in \Omega} |f(x)|$ .

**Proposition 18.0.10.** Let  $\Omega$  be a closed bounded subset of  $\mathbb{R}^n$ . The uniform convergence norm is a norm on  $\mathcal{C}(\Omega)$ .

$$\begin{split} \langle 1 \rangle 1. & \text{ If } \|f\| = 0 \text{ then } f = 0 \\ & \text{Proof: If } \max_x |f(x)| = 0 \text{ then } f(x) = 0 \text{ for all } x. \\ \langle 1 \rangle 2. & \|\lambda f\| = |\lambda| \|f\| \\ & \text{Proof:} \\ & \|\lambda f\| = \max_x |\lambda f(x)| \\ & = |\lambda| \max_x |f(x)| \\ & = |\lambda| \|f\| \end{split}$$

 $\langle 1 \rangle 3. \| f + g \| \le \| f \| + \| g \|$ PROOF:

$$||f + g|| = \max_{x} |f(x) + g(x)|$$

$$\leq \max_{x} (|f(x)| + |g(x)|)$$

$$\leq \max_{x} |f(x)| + \max_{x} |g(x)|$$

$$= ||f|| + ||g||$$

**Proposition 18.0.11.** Let  $p \ge 1$ . The function  $||(z_n)|| = (\sum_{n=1}^{\infty} |z_n|^p)^{1/p}$  is a norm on  $l^p$ .

Proof:

 $\langle 1 \rangle 1$ . If  $||(z_n)|| = 0$  then  $(z_n) = (0)$ PROOF: If  $(\sum_n |z_n|^p)^{1/p} = 0$  then  $\sum_n |z_n|^p = 0$  so  $|z_n|^p = 0$  for all n, and so  $z_n = 0$  for all n.

 $\langle 1 \rangle 2$ .  $\|(\lambda z_n)\| = |\lambda| \|(z_n)\|$ 

Proof:

$$\|(\lambda z_n)\| = \left(\sum_n |\lambda z_n|^p\right)^{1/p}$$
$$= |\lambda| \left(\sum_n |z_n|^p\right)^{1/p}$$
$$= |\lambda| |(z_n)|$$

 $\langle 1 \rangle 3$ . The triangle inequality holds.

PROOF: This is Minkowski's Inequality.

**Proposition 18.0.12.** Let V be a normed space and U a vector subspace of V. Then U is a normed space under the restriction of the norm to U.

PROOF: Each axiom follows from the fact it holds in V.  $\square$ 

**Proposition 18.0.13.** Let V be a normed space over K. Let  $x_1, \ldots, x_n$  be linearly independent elements of V. Then there exists a real number c > 0 such that, for all  $\alpha_1, \ldots, \alpha_n \in K$ , we have

$$\|\alpha_1 x_1 + \dots + \alpha_n x_n\| \ge c(|\alpha_1| + \dots + |\alpha_n|).$$

Proof:

 $\langle 1 \rangle 1$ . Define  $f: K^n \to \mathbb{R}$  by

$$f(\alpha_1, \dots, \alpha_n) = \|\alpha_1 x_1 + \dots + \alpha_n x_n\|$$

 $\langle 1 \rangle 2$ . f is continuous.

 $\langle 2 \rangle 1$ . Let:  $(\alpha_1, \ldots, \alpha_n) \in K^n$  and  $\epsilon > 0$ 

 $\langle 2 \rangle 2$ . Let:  $\delta = \epsilon/(\|x_1\| + \cdots + \|x_n\|)$ 

PROOF:  $x_1, \ldots, x_n$  are not all zero because they are linearly independent.

 $\langle 2 \rangle 3$ . Let:  $(\beta_1, \ldots, \beta_n)$  with  $|\alpha_i - \beta_i| < \delta$  for all i

```
\langle 2 \rangle 4. \|(\alpha_1 x_1 + \dots + \alpha_n x_n) - (\beta_1 x_1 + \beta_n x_n)\| < \epsilon
                  \|(\alpha_1 x_1 + \dots + \alpha_n x_n) - (\beta_1 x_1 + \beta_n x_n)\|
               \leq |\alpha_1 - \beta_1| ||x_1|| + \dots + |\alpha_n - \beta_n| ||x_n||
                                                                                      (Axioms 2 and 3)
               <\delta(||x_1|| + \cdots + ||x_n||)
                                                                                                         (\langle 2 \rangle 3)
                                                                                                         (\langle 2 \rangle 2)
\langle 1 \rangle 3. Pick (\beta_1, \dots, \beta_n) \in \{(\beta_1, \dots, \beta_n) \in K^n \mid |\beta_1| + \dots + |\beta_n| = 1\} at which
         f attains its minimum.
   PROOF: Extreme Value Theorem.
\langle 1 \rangle 4. Let c = f(\beta_1, \dots, \beta_n)
\langle 1 \rangle 5. \ c > 0
   Proof: Linear independence.
\langle 1 \rangle 6. Let: \alpha_1, \ldots, \alpha_n \in K
\langle 1 \rangle 7. \|\alpha_1 x_1 + \dots + \alpha_n x_n\| \ge c(|\alpha_1| + \dots + |\alpha_n|)
   \langle 2 \rangle 1. Assume: w.l.o.g. \alpha_1 \ldots, \alpha_n are not all zero.
   \langle 2 \rangle 2. Let: \beta_i = \alpha_i/(|\alpha_1| + \cdots + |\alpha_n|) for i = 1, \dots, n
   \langle 2 \rangle 3. |\beta_1| + \cdots + |\beta_n| = 1
   \langle 2 \rangle 4. \ f(\beta_1, \dots, \beta_n) \ge c
   \langle 2 \rangle5. Q.E.D.
       PROOF: Multiply both sides by |\alpha_1| + \cdots + |\alpha_n|.
П
Proposition 18.0.14. Let V be a normed space over K. Define d: V^2 \to \mathbb{R}
by d(x,y) = ||x-y||. Then d is a metric on V.
Proof:
\langle 1 \rangle 1. For all x, y \in V we have d(x, y) \geq 0
   Proof: Proposition 18.0.3.
\langle 1 \rangle 2. For all x, y \in V we have d(x, y) = 0 iff x = y
   \langle 2 \rangle 1. If d(x,y) = 0 then x = y
       Proof: Axiom 1.
   \langle 2 \rangle 2. If x = y then d(x, y) = 0
       Proof: Proposition 18.0.2.
\langle 1 \rangle 3. \ \forall x, y \in V.d(x, y) = d(y, x)
   PROOF: By Axiom 2.
\langle 1 \rangle 4. \ \forall x, y, z \in V.d(x, z) \le d(x, y) + d(y, z)
   Proof: By Axiom 3.
```

Henceforth we identify any normed space with this metric space.

### 18.1 Convergence

**Proposition 18.1.1.** Let V be a normed space over K. Let  $(x_n)$  be a sequence in V and  $l \in V$ . Then  $x_n \to l$  as  $n \to \infty$  in V if and only if  $||x_n - l|| \to 0$  as  $n \to \infty$  in  $\mathbb{R}$ .

PROOF: Immediate from definitions.  $\square$ 

Proposition 18.1.2. In a normed space, a sequence has at most one limit.

#### Proof:

- $\langle 1 \rangle 1$ . Let: V be a vector space over K.
- $\langle 1 \rangle 2$ . Assume:  $x_n \to l$  and  $x_n \to m$  as  $n \to \infty$ .
- $\langle 1 \rangle 3$ . Assume: for a contradiction  $l \neq m$
- $\langle 1 \rangle 4$ . Let:  $\epsilon = ||l m||/2$
- $\langle 1 \rangle$ 5. PICK N such that  $\forall n \geq N. ||x_n l|| < \epsilon$  and  $\forall n \geq N. ||x_n m|| < \epsilon$  PROOF:  $\langle 1 \rangle 2$ ,  $\langle 1 \rangle 4$
- $\langle 1 \rangle 6. \ \|l m\| < \|l m\|$

Proof:

$$\begin{split} \|l-m\| &\leq \|x_N-l\| + \|x_N-m\| & \text{(Triangle Inequality)} \\ &< 2\epsilon & \text{($\langle 1\rangle 5$)} \\ &= \|l-m\| & \text{($\langle 1\rangle 4$)} \end{split}$$

 $\langle 1 \rangle 7$ . Q.E.D.

Proof: This is a contradiction.

**Definition 18.1.3** (Bounded). Let V be a normed space over K. A sequence  $(x_n)$  in V is bounded iff there exists B such that  $\forall n \leq N . ||x_n|| < B$ .

Proposition 18.1.4. Every convergent sequence is bounded.

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $x_n \to l$  as  $n \to \infty$
- $\langle 1 \rangle 2$ . PICK N such that  $\forall n \geq N . ||x_n l|| < 1$
- $\langle 1 \rangle 3$ . Let:  $B = \max(||x_1||, ||x_2||, \dots, ||x_{N-1}||, ||l|| + 1)$
- $\langle 1 \rangle 4$ . Let:  $n \in \mathbb{N}$
- $\langle 1 \rangle 5. \|x_n\| \leq B$ 
  - $\langle 2 \rangle 1$ . Case: n < N

PROOF:  $||x_n|| \leq B$  from  $\langle 1 \rangle 3$ .

 $\langle 2 \rangle 2$ . Case:  $n \geq N$ 

Proof:

$$||x_n|| \le ||l|| + ||x_n - l||$$
 (Triangle Inequality)  
 $< ||l|| + 1$  ( $\langle 1 \rangle 2$ )  
 $\le B$  ( $\langle 1 \rangle 3$ )

**Proposition 18.1.5.** Let V be a normed space over K. If  $x_n \to l$  as  $n \to \infty$  in V, and  $\lambda_n \to \lambda$  as  $n \to \infty$  in K, then  $\lambda_n x_n \to \lambda l$  as  $n \to \infty$ .

- $\langle 1 \rangle 1$ . Let: V be a normed space over K.
- $\langle 1 \rangle 2$ . Let:  $x_n \to l$  as  $n \to \infty$
- $\langle 1 \rangle 3$ . Let:  $\lambda_n \to \lambda$  as  $n \to \infty$

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$$\begin{array}{ll} \langle 1 \rangle 4. & \text{Let: } \epsilon > 0 \\ \langle 1 \rangle 5. & \text{Pick } N \text{ such that, for all } n \geq N, \text{ we have } \|x_n - l\| < \epsilon/2 |\lambda| \text{ and } |\lambda_n - \lambda| < \sqrt{\epsilon/2} \text{ and } \|x_n\| < \sqrt{\epsilon/2} \\ \langle 1 \rangle 6. & \text{Let: } n \geq N \\ \langle 1 \rangle 7. & \|\lambda_n x_n - \lambda l\| < \epsilon \\ & \text{Proof:} \\ & \|\lambda_n x_n - \lambda l\| \leq \|\lambda_n x_n - \lambda x_n\| + \|\lambda x_n - \lambda l\| & \text{(Triangle Inequality)} \\ & = |\lambda_n - \lambda| \|x_n\| + |\lambda| \|x_n - l\| & \text{(Axiom 2)} \\ & < \sqrt{\epsilon/2} \sqrt{\epsilon/2} + |\lambda| \epsilon/2 |\lambda| & \text{($\langle 1 \rangle 5$)} \\ & = \epsilon \end{array}$$

**Proposition 18.1.6.** Let V be a normed space over K. If  $x_n \to l$  and  $y_n \to m$  as  $n \to \infty$ , then  $x_n + y_n \to l + m$  as  $n \to \infty$ .

Proof:

$$\langle 1 \rangle 1$$
. Let:  $\epsilon > 0$ 

 $\langle 1 \rangle 2$ . PICK N such that, for all  $n \geq N$ , we have  $||x_n - l|| < \epsilon/2$  and  $||y_n - m|| < \epsilon/2$ 

$$\langle 1 \rangle 3$$
. Let:  $n \geq N$ 

$$\langle 1 \rangle 4. \ \|(x_n + y_n) - (l+m)\| < \epsilon$$

Proof:

$$\begin{aligned} \|(x_n+y_n)-(l+m)\| &\leq \|x_n-l\|+\|y_n-m\| &\quad \text{(Triangle Inequality)} \\ &<\epsilon/2+\epsilon/2 &\quad (\langle 1\rangle 2) \\ &=\epsilon \end{aligned}$$

**Definition 18.1.7** (Uniform Convergence). Let  $\Omega$  be a closed bounded subset of  $\mathbb{R}^n$ . Let  $(f_n)$  be a sequence in  $\mathcal{C}(\Omega)$  and  $f \in \mathcal{C}(\Omega)$ . Then  $(f_n)$  converges uniformly to f iff, for every  $\epsilon > 0$ , there exists N such that  $\forall x \in \Omega. \forall n \geq N. |f_n(x) - f(x)| < \epsilon$ .

**Proposition 18.1.8.** Let  $\Omega$  be a closed bounded subset of  $\mathbb{R}^n$ . Let  $(f_n)$  be a sequence in  $\mathcal{C}(\Omega)$  and  $f \in \mathcal{C}(\Omega)$ . Then  $(f_n)$  converges uniformly to f iff  $f_n$  converges to f under the uniform convergence norm.

Proof:

$$(f_n)$$
 converges to  $f$  under the uniform convergence norm  $\Leftrightarrow \forall \epsilon > 0. \exists N. \forall n \geq N. ||f_n - f|| < \epsilon$   $\Leftrightarrow \forall \epsilon > 0. \exists N. \forall n \geq N. \forall x \in X. |f_n(x) - f(x)| < \epsilon$ 

**Definition 18.1.9** (Pointwise Convergence). Let  $(f_n)$  be a sequence in  $\mathcal{C}([0,1])$  and  $f \in \mathcal{C}([0,1])$ . Then  $(f_n)$  converges pointwise to f iff, for all  $t \in [0,1]$ , we have  $|f_n(t) - f(t)| \to 0$  as  $n \to \infty$ .

**Proposition 18.1.10.** There is no norm n on C([0,1]) such that, for every sequence  $(f_n)$  and function f in C([0,1]), we have  $(f_n)$  converges pointwise to f if and only if  $(f_n)$  converges to f under n.

Proof:

 $\langle 1 \rangle 1$ . Assume: for a contradiction  $\| \|$  is a norm on  $\mathcal{C}([0,1])$  such that, for every sequence  $(f_n)$  and function f in  $\mathcal{C}([0,1])$ , we have  $(f_n)$  converges pointwise to f if and only if  $(f_n)$  converges to f under  $\| \|$ .

 $\langle 1 \rangle 2$ . For  $n \in \mathbb{Z}_+$ , define  $g_n \in \mathcal{C}([0,1])$  by

$$g_n(t) = \begin{cases} 2^n t & \text{if } 0 \le t \le 2^{-n} \\ 2 - 2^n t & \text{if } 2^{-n} \le t \le 2^{1-n} \\ 0 & \text{if } 2^{1-n} \le t \le 1 \end{cases}$$

 $\langle 1 \rangle 3$ . For all n,  $||g_n|| \neq 0$ 

Proof: Axiom 1.

 $\langle 1 \rangle 4$ . For  $n \in \mathbb{Z}_+$ , define  $f_n \in \mathcal{C}([0,1])$  by  $f_n = g_n/\|g_n\|$ 

 $\langle 1 \rangle 5$ . For all n,  $||f_n|| = 1$ 

Proof: Axiom 2.

 $\langle 1 \rangle 6$ .  $(f_n)$  does not converge under  $\| \|$ 

 $\langle 1 \rangle 7$ .  $(f_n)$  converges pointwise to 0.

 $\langle 1 \rangle 8$ . This is a contradiction.

**Definition 18.1.11** (Equivalence of Norms). Let  $\| \|_1$  and  $\| \|_2$  be two norms on the same vector space V. Then the norms are *equivalent* if and only if, for any sequence  $(x_n)$  in V and  $l \in V$ , we have that  $(x_n)$  converges to l under  $\| \|_1$  if and only if  $(x_n)$  converges to l under  $\| \|_2$ .

**Theorem 18.1.12.** Let  $\| \|_1$  and  $\| \|_2$  be two norms on the same vector space E over K. Then  $\| \|_1$  and  $\| \|_2$  are equivalent if and only if there exist positive real numbers  $\alpha$  and  $\beta$  such that, for all  $x \in E$ ,

$$\alpha ||x||_1 \le ||x||_2 \le \beta ||x||_1$$
.

- $\langle 1 \rangle 1$ . If  $\| \|_1$  and  $\| \|_2$  are equivalent then there exist positive real numbers  $\alpha$  and  $\beta$  such that, for all  $x \in E$ ,  $\alpha \|x\|_1 \le \|x\|_2 \le \beta \|x\|_1$ .
  - $\langle 2 \rangle 1$ . Assume:  $\| \|_1$  and  $\| \|_2$  are equivalent.
  - $\langle 2 \rangle 2$ . There exists  $\alpha > 0$  such that, for all  $x \in E$ , we have  $\alpha ||x||_1 \leq ||x||_2$ 
    - $\langle 3 \rangle 1$ . Assume: for a contradiction there is no  $\alpha > 0$  such that, for all  $x \in E$ , we have  $\alpha ||x||_1 \le ||x||_2$ .
    - $\langle 3 \rangle 2$ . For all  $n \in \mathbb{Z}_+$ , PICK  $x_n \in E$  such that  $1/n ||x_n||_1 > ||x||_2$
    - $\langle 3 \rangle 3$ . For all  $n \in \mathbb{Z}_+$ , Let:

$$y_n = \frac{1}{\sqrt{n}} \frac{x_n}{\|x_n\|_2}$$

- $\langle 3 \rangle 4$ .  $(y_n)$  converges to 0 under  $\| \|_2$
- $\langle 3 \rangle 5.$   $(y_n)$  converges to 0 under  $\| \|_1$
- $\langle 3 \rangle 6$ . For all  $n \in \mathbb{Z}_+$ , we have  $||y_n|| > \sqrt{n}$
- $\langle 3 \rangle 7$ . This is a contradiction.
- $\langle 2 \rangle$ 3. There exists  $\beta > 0$  such that, for all  $x \in E$ , we have  $||x||_2 \le \beta ||x||_1$  PROOF: Similar.

- $\langle 1 \rangle 2$ . If there exist positive real numbers  $\alpha$  and  $\beta$  such that, for all  $x \in E$ ,  $\alpha \|x\|_1 \leq \|x\|_2 \leq \beta \|x\|_1$ , then  $\| \ \|_1$  and  $\| \ \|_2$  are equivalent.
  - $\langle 2 \rangle 1$ . Assume:  $\alpha$  and  $\beta$  are positive reals with  $\forall x \in E.\alpha ||x||_1 \leq ||x||_2 \leq \beta ||x||_1$ .
  - $\langle 2 \rangle 2$ . Let  $(x_n)$  be a sequence in E and  $l \in E$
  - $\langle 2 \rangle 3$ . If  $(x_n)$  converges to l under  $\| \|_1$  then  $(x_n)$  converges to l under  $\| \|_2$ .
    - $\langle 3 \rangle 1$ . Assume:  $(x_n)$  converges to l under  $\| \|_1$
    - $\langle 3 \rangle 2$ . Let:  $\epsilon > 0$
    - $\langle 3 \rangle 3$ . PICK N such that  $\forall n \geq N . ||x_n l||_1 < \epsilon/\beta$
    - $\langle 3 \rangle 4. \ \forall n \geq N. ||x_n l||_2 < \epsilon$
  - $\langle 2 \rangle 4$ . If  $(x_n)$  converges to l under  $\| \|_2$  then  $(x_n)$  converges to l under  $\| \|_1$ . PROOF: Similar.

**Theorem 18.1.13.** Any two norms on a finite dimensional vector space are equivalent.

#### Proof:

- $\langle 1 \rangle 1$ . Let: V be a finite dimensional vector space over K.
- $\langle 1 \rangle 2$ . Assume: w.l.o.g. dim V > 0
- $\langle 1 \rangle 3$ . PICK a basis  $\{e_1, \ldots, e_n\}$  for V.
- $\langle 1 \rangle 4$ . Let:  $\| \|_0 : V \to \mathbb{R}$  be the function:  $\| \alpha_1 e_1 + \dots + \alpha_n e_n \|_0 = |\alpha_1| + \dots + |\alpha_n|$ .
- $\langle 1 \rangle 5$ .  $\| \|_0$  is a norm.
  - $\langle 2 \rangle 1$ . If  $||v||_0 = 0$  then v = 0

PROOF: If  $|\alpha_1| + \dots + |\alpha_n| = 0$  then  $\alpha_1 = \dots = \alpha_n = 0$  so  $\alpha_1 e_1 + \dots + \alpha_n e_n = 0$ 

 $\langle 2 \rangle 2$ .  $\|\lambda v\|_0 = |\lambda| \|v\|_0$ 

Proof:

$$\|\lambda(\alpha_1 e_1 + \dots + \alpha_n e_n)\|_0 = \|\lambda \alpha_1 e_1 + \dots + \lambda \alpha_n e_n\|_0$$

$$= |\lambda \alpha_1| + \dots + |\lambda \alpha_n| \qquad (\langle 1 \rangle 4)$$

$$= |\lambda|(|\alpha_1| + \dots + |\alpha_n|)$$

$$= |\lambda|\|\alpha_1 e_1 + \dots + \alpha_n e_n\|_0 \qquad (\langle 1 \rangle 4)$$

 $\langle 2 \rangle 3. \|u + v\|_0 \le \|u\|_0 + \|v\|_0$ 

Proof:

$$\|(\alpha_1 e_1 + \dots + \alpha_n e_n) + (\beta_1 e_1 + \dots + \beta_n e_n)\| = |\alpha_1 + \beta_1| + \dots + |\alpha_n + \beta_n|$$

$$\leq |\alpha_1| + \dots + |\alpha_n| + |\beta_1| + \dots + |\beta_n|$$

$$= \|\alpha_1 e_1 + \dots + \alpha_n e_n\|_0 + \|\beta_1 e_1 + \dots + \beta_n e_n\|_0$$

- $\langle 1 \rangle 6$ . Any norm on V is equivalent to  $\| \|_0$ .
  - $\langle 2 \rangle 1$ . Let:  $\| \|$  be any norm on V.
  - $\langle 2 \rangle 2$ . PICK  $\alpha > 0$  such that, for all  $\alpha_1, \ldots, \alpha_n \in K$ , we have  $\|\alpha_1 e_1 + \cdots + \alpha_n e_n\| \ge \alpha(|\alpha_1| + \cdots + |\alpha_n|)$

PROOF: Proposition 18.0.13,  $\langle 2 \rangle 1$ ,  $\langle 1 \rangle 3$ .

- $\langle 2 \rangle 3$ . Let:  $\beta = \max(\|e_1\|, \dots, \|e_n\|)$
- $\langle 2 \rangle 4. \ \beta > 0$

PROOF:  $e_1, \ldots, e_n$  cannot all be zero by  $\langle 1 \rangle 3$ .

- $\langle 2 \rangle 5$ . For all  $x \in V$  we have  $\alpha ||x||_0 \le ||x|| \le \beta ||x||_0$ 
  - $\langle 3 \rangle 1$ . Let:  $x \in V$
  - $\langle 3 \rangle 2$ .  $\alpha ||x||_0 \leq ||x||$

Proof:  $\langle 1 \rangle 3$ ,  $\langle 1 \rangle 4$ ,  $\langle 2 \rangle 2$ .

 $\langle 3 \rangle 3$ .  $||x|| \leq \beta ||x||_0$ 

 $\langle 4 \rangle 1$ . Let:  $x = \alpha_1 e_1 + \cdots + \alpha_n e_n$ 

 $\langle 4 \rangle 2$ . Q.E.D.

Proof:

$$||x|| = ||\alpha_1 e_1 + \dots + \alpha_n e_n|| \qquad (\langle 4 \rangle 1)$$

$$\leq |\alpha_1| ||e_1|| + \dots + |\alpha_n| ||e_n|| \qquad (\langle 2 \rangle 1)$$

$$\leq \beta(|\alpha_1| + \dots + |\alpha_n|) \tag{(2)3}$$

$$=\beta \|x\|_0 \tag{(1)4}$$

 $\langle 2 \rangle 6$ . Q.E.D.

PROOF: Theorem 18.1.12,  $\langle 1 \rangle 5$ ,  $\langle 2 \rangle 1$ ,  $\langle 2 \rangle 2$ ,  $\langle 2 \rangle 4$ ,  $\langle 2 \rangle 5$ .

**Definition 18.1.14** (Open Ball). Let V be a normed space over K. Let  $x \in V$ . Let r > 0. The *open ball* with *centre* x and *radius* r is

$$B(x,r) := \{ y \in V \mid ||y - x|| < r \} .$$

**Definition 18.1.15** (Closed Ball). Let V be a normed space over K. Let  $x \in V$ . Let r > 0. The *closed ball* with *centre* x and *radius* r is

$$\overline{B}(x,r) := \{ y \in V \mid ||y - x|| \le r \}$$
.

**Definition 18.1.16** (Sphere). Let V be a normed space over K. Let  $x \in V$ . Let r > 0. The *sphere* with *centre* x and *radius* r is

$$S(x,r) := \{ y \in V \mid ||y - x|| = r \} .$$

**Definition 18.1.17** (Open Set). Let V be a normed space over K. A set  $S \subseteq V$  is *open* iff, for all  $x \in S$ , there exists  $\epsilon > 0$  such that  $B(x, \epsilon) \subseteq S$ .

**Proposition 18.1.18.** Equivalent norms define the same set of open sets.

- $\langle 1 \rangle 1$ . Let: V be a normed space over K.
- $\langle 1 \rangle 2$ . Let:  $\| \|_1$  and  $\| \|_2$  be equivalent norms on V.
- (1)3. PICK reals  $\alpha, \beta > 0$  such that, for all  $x \in V$ , we have  $\alpha ||x||_1 \le ||x||_2 \le \beta ||x||_1$
- $\langle 1 \rangle 4$ . Let:  $S \subseteq V$
- $\langle 1 \rangle 5$ . If S is open under  $\| \|_1$  then S is open under  $\| \|_2$ .
  - $\langle 2 \rangle 1$ . Assume: S is open under  $\| \|_1$ .
  - $\langle 2 \rangle 2$ . Let:  $x \in S$
  - $\langle 2 \rangle 3$ . Pick  $\epsilon > 0$  such that  $\{ y \in V \mid ||x y||_1 < \epsilon \} \subseteq S$ .
  - $\langle 2 \rangle 4$ . Let:  $\delta = \alpha \epsilon$

```
\langle 2 \rangle5. \{ y \in V \mid \|x - y\|_2 < \delta \} \subseteq S
\langle 1 \rangle6. If S is open under \| \ \|_2 then S is open under \| \ \|_1.
PROOF: Similar.
```

Proposition 18.1.19. Every open ball is open.

PROOF:

 $\langle 1 \rangle 1$ . Let: V be a normed space over K.

 $\langle 1 \rangle 2$ . Let:  $c \in V$  and r > 0Prove: B(c, r) is open.

 $\langle 1 \rangle 3$ . Let:  $x \in B(c,r)$ 

 $\langle 1 \rangle 4$ . Let:  $\epsilon = r - ||x - c||$ Prove:  $B(x, \epsilon) \subseteq B(c, r)$ 

 $\langle 1 \rangle$ 5. Let:  $y \in B(x, \epsilon)$ Prove:  $y \in B(c, r)$ 

 $\langle 1 \rangle 6$ . ||y - c|| < r

Proof:

$$\begin{split} \|y-c\| &\leq \|y-x\| + \|x-c\| & \text{(Triangle Inequality)} \\ &< \epsilon + \|x-c\| & \text{($\langle 1 \rangle 5$)} \\ &= r & \text{($\langle 1 \rangle 4$)} \end{split}$$

**Proposition 18.1.20.** Let  $\Omega$  be a closed bounded set in  $\mathbb{R}^n$ . Let  $f \in \mathcal{C}(\Omega)$ . Then  $U = \{g \in \mathcal{C}(\Omega) \mid \forall x \in \Omega. g(x) < f(x)\}$  is open.

Proof:

 $\langle 1 \rangle 1$ . Let:  $g \in U$ 

 $\langle 1 \rangle 2$ . Let:  $\epsilon = \max_{x \in \Omega} (f(x) - g(x))$ Prove:  $B(g, \epsilon) \subseteq S$ 

 $\langle 1 \rangle 3. \ \epsilon > 0$ 

 $\langle 1 \rangle 4$ . Let:  $h \in B(g, \epsilon/2)$ 

Prove:  $h \in S$ 

 $\langle 1 \rangle 5$ . Let:  $x \in \Omega$ 

 $\langle 1 \rangle 6. \ h(x) < f(x)$ 

Proof:

$$h(x) \le g(x) + \epsilon/2$$
  $(\langle 1 \rangle 4)$ 

$$\langle g(x) + \epsilon \rangle$$
 ( $\langle 1 \rangle 3$ )

$$\leq f(x)$$
  $(\langle 1 \rangle 2)$ 

**Proposition 18.1.21.** Let  $\Omega$  be a closed bounded set in  $\mathbb{R}^n$ . Let  $f \in \mathcal{C}(\Omega)$ . Then  $U = \{g \in \mathcal{C}(\Omega) \mid \forall x \in \Omega. g(x) > f(x)\}$  is open.

PROOF: Given  $g \in U$ , let  $\epsilon = \max_x (g(x) - f(x))/2$ . Then  $B(g, \epsilon) \subseteq U$ .  $\square$ 

**Proposition 18.1.22.** Let  $\Omega$  be a closed bounded set in  $\mathbb{R}^n$ . Let  $f \in \mathcal{C}(\Omega)$  be such that f(x) > 0 for all  $x \in \Omega$ . Then  $U = \{g \in \mathcal{C}(\Omega) \mid \forall x \in \Omega. |g(x)| < f(x)\}$  is open.

PROOF: Given  $g \in U$ , let  $\epsilon = \max_x (f(x) - |g(x)|)/2$ . Then  $B(g, \epsilon) \subseteq U$ .  $\square$ 

**Proposition 18.1.23.** Let  $\Omega$  be a closed bounded set in  $\mathbb{R}^n$ . Let  $f \in \mathcal{C}(\Omega)$  be such that f(x) > 0 for all  $x \in \Omega$ . Then  $U = \{g \in \mathcal{C}(\Omega) \mid \forall x \in \Omega. |g(x)| > f(x)\}$  is open.

PROOF: Given  $g \in U$ , let  $\epsilon = \max_{x} (|g(x)| - f(x))/2$ . Then  $B(g, \epsilon) \subseteq U$ .  $\square$ 

Proposition 18.1.24. The union of a set of open sets is open.

#### Proof:

- $\langle 1 \rangle 1$ . Let: V be a normed space over K.
- $\langle 1 \rangle 2$ . Let:  $\mathcal{U}$  be a set of open sets in V.
- $\langle 1 \rangle 3$ . Let:  $x \in \bigcup \mathcal{U}$
- $\langle 1 \rangle 4$ . PICK  $U \in \mathcal{U}$  such that  $x \in U$ .
- $\langle 1 \rangle 5$ . Pick  $\epsilon > 0$  such that  $B(x, \epsilon) \subseteq U$
- $\langle 1 \rangle 6. \ B(x, \epsilon) \subseteq \bigcup \mathcal{U}$

Proposition 18.1.25. The intersection of two open sets is open.

#### **PROOF**

- $\langle 1 \rangle 1$ . Let: V be a normed space over K.
- $\langle 1 \rangle 2$ . Let:  $U_1$  and  $U_2$  be open sets in V.
- $\langle 1 \rangle 3$ . Let:  $x \in U_1 \cap U_2$
- $\langle 1 \rangle 4$ . Pick  $\epsilon_1 > 0$  such that  $B(x, \epsilon_1) \subseteq U_1$
- $\langle 1 \rangle 5$ . Pick  $\epsilon_2 > 0$  such that  $B(x, \epsilon_2) \subseteq U_2$
- $\langle 1 \rangle 6$ . Let:  $\epsilon = \min(\epsilon_1, \epsilon_2)$
- $\langle 1 \rangle 7. \ B(x,\epsilon) \subseteq U_1 \cap U_2$

**Proposition 18.1.26.** In any normed space,  $\emptyset$  is open.

Proof: Vacuous.

**Proposition 18.1.27.** In any normed space V, the whole space V is open.

PROOF: For any  $x \in V$  we have  $B(x,1) \subseteq V$ .  $\square$ 

**Definition 18.1.28** (Closed Set). Let V be a normed space over K. A set  $S \subseteq V$  is *closed* iff V - S is open.

Proposition 18.1.29. Every closed ball is closed.

- $\langle 1 \rangle 1$ . Let: V be a normed space over K.
- $\begin{array}{ll} \langle 1 \rangle 2. \ \ \mathrm{Let:} \ \ c \in V \ \ \mathrm{and} \ \ r > 0 \\ & \mathrm{Prove:} \ \ \overline{B}(c,r) \ \mathrm{is} \ \mathrm{closed}. \end{array}$
- $\langle 1 \rangle 3$ . Let:  $x \in V \overline{B}(c, r)$
- $\langle 1 \rangle 4$ . Let:  $\epsilon = ||x c|| r$ Prove:  $B(x, \epsilon) \subseteq V - \overline{B}(c, r)$

П

$$\begin{split} \langle 1 \rangle 5. & \epsilon > 0 \\ \text{PROOF: Since } \|x - c\| > r \text{ by } \langle 1 \rangle 3. \\ \langle 1 \rangle 6. \text{ Let: } y \in B(x, \epsilon) \\ \langle 1 \rangle 7. & \|y - c\| > r \\ \text{PROOF:} \\ & \|y - c\| \geq \|x - c\| - \|x - y\| \\ & > \|x - c\| - \epsilon \\ & = r \end{split} \tag{Triangle Inequality}$$

Proposition 18.1.30. The intersection of a set of closed sets is closed.

PROOF: From Proposition 18.1.24.

Proposition 18.1.31. The union of two closed sets is closed.

Proof: From Proposition 18.1.25.  $\square$ 

Proposition 18.1.32. Every sphere is closed.

PROOF:  $S(c,r) = \overline{B}(c,r) - B(c,r)$ .

**Proposition 18.1.33.** Let  $\Omega$  be a closed bounded set in  $\mathbb{R}^n$ . Let  $f \in \mathcal{C}(\Omega)$ . Then  $\{g \in \mathcal{C}(\Omega) \mid \forall x \in \Omega. g(x) \leq f(x)\}$  is closed.

PROOF: It is  $\mathcal{C}(\Omega) - \{g \mid \forall x \in \Omega. g(x) > f(x)\}.$ 

**Proposition 18.1.34.** Let  $\Omega$  be a closed bounded set in  $\mathbb{R}^n$ . Let  $f \in \mathcal{C}(\Omega)$ . Then  $\{g \in \mathcal{C}(\Omega) \mid \forall x \in \Omega. g(x) \geq f(x)\}$  is closed.

PROOF: It is  $\mathcal{C}(\Omega) - \{g \mid \forall x \in \Omega. g(x) < f(x)\}.$ 

**Proposition 18.1.35.** Let  $\Omega$  be a closed bounded set in  $\mathbb{R}^n$ . Let  $f \in \mathcal{C}(\Omega)$ . Then  $\{g \in \mathcal{C}(\Omega) \mid \forall x \in \Omega. | g(x)| \leq f(x)\}$  is closed.

PROOF: It is  $\mathcal{C}(\Omega) - \{g \mid \forall x \in \Omega . |g(x)| > f(x)\}.$ 

**Proposition 18.1.36.** Let  $\Omega$  be a closed bounded set in  $\mathbb{R}^n$ . Let  $f \in \mathcal{C}(\Omega)$ . Then  $\{g \in \mathcal{C}(\Omega) \mid \forall x \in \Omega. | g(x)| \geq f(x)\}$  is closed.

PROOF: It is  $C(\Omega) - \{g \mid \forall x \in \Omega . |g(x)| < f(x)\}.$ 

**Proposition 18.1.37.** Let  $\Omega$  be a closed bounded set in  $\mathbb{R}^n$ . Let  $x_0 \in \Omega$  and  $\lambda \in \mathbb{C}$ . Then  $C = \{g \in \mathcal{C}(\Omega) \mid g(x_0) = \lambda\}$  is closed.

PROOF: Given  $g \in \mathcal{C}(\Omega) - C$ , let  $\epsilon = |g(x_0) - \lambda|/2$ . Then  $B(g, \epsilon) \subseteq \mathcal{C}(\Omega) - C$ .  $\square$ 

**Proposition 18.1.38.** In any normed space V, we have  $\emptyset$  is closed.

PROOF: Since  $V - \emptyset = V$  is open.  $\square$ 

**Proposition 18.1.39.** In any normed space V, the whole space V is closed.

PROOF: Since  $V - V = \emptyset$  is open.  $\square$ 

**Theorem 18.1.40.** Let V be a normed space over K. Let S be a subset of V. Then S is closed if and only if, for any sequence  $(x_n)$  in S, if  $x_n \to l$  as  $n \to \infty$  then  $l \in S$ .

#### Proof:

- $\langle 1 \rangle 1$ . If S is closed then, for any sequence  $(x_n)$  in S, if  $x_n \to l$  as  $n \to \infty$  then  $l \in S$ .
  - $\langle 2 \rangle 1$ . Assume: S is closed.
  - $\langle 2 \rangle 2$ . Let:  $(x_n)$  be a sequence in S.
  - $\langle 2 \rangle 3$ . Assume:  $x_n \to l$  as  $n \to \infty$ .
  - $\langle 2 \rangle 4$ . Assume: for a contradiction  $l \notin S$ .
  - $\langle 2 \rangle$ 5. Pick  $\epsilon > 0$  such that  $B(l, \epsilon) \subseteq V S$
  - $\langle 2 \rangle 6$ . Pick N such that  $\forall n \geq N.x_n \in B(l, \epsilon)$
  - $\langle 2 \rangle 7. \ x_N \in V S$
  - $\langle 2 \rangle 8$ . This contradicts  $\langle 2 \rangle 2$ .
- $\langle 1 \rangle 2$ . If, for any sequence  $(x_n)$  in S, if  $x_n \to l$  as  $n \to \infty$  then  $l \in S$ , then S is closed.
  - $\langle 2 \rangle 1$ . Assume: for any sequence  $(x_n)$  in S, if  $x_n \to l$  as  $n \to \infty$  then  $l \in S$ .
  - $\langle 2 \rangle 2$ . Let:  $x \in V S$
  - $\langle 2 \rangle 3$ . Assume: for a contradiction there is no  $\epsilon > 0$  such that  $B(x, \epsilon) \subseteq V S$ .
  - $\langle 2 \rangle 4$ . For  $n \in \mathbb{Z}_+$ , Pick  $x_n \in B(x, 1/n) \cap S$
  - $\langle 2 \rangle 5. \ x_n \to x \text{ as } n \to \infty$
  - $\langle 2 \rangle 6. \ x \in S$
  - $\langle 2 \rangle 7$ . This contradicts  $\langle 2 \rangle 2$ .

**Definition 18.1.41** (Closure). Let V be a normed space over K. Let S be a subset of V. The *closure* of S,  $\operatorname{cl} S$ , is the intersection of the set of closed sets that include S.

**Proposition 18.1.42.** Let V be a normed space over K. Let S be a subset of V. Then the closure of S is the smallest closed set that includes S.

Proof: Proposition 18.1.30.

**Theorem 18.1.43.** Let V be a normed space over K. Let S be a subset of V. Then

$$\operatorname{cl} S = \{ l \in V \mid \exists \text{ a sequence } (x_n) \text{ in } S.x_n \to l \text{ as } n \to \infty \} .$$

#### Proof:

- $\langle 1 \rangle 1$ . For all  $l \in \operatorname{cl} S$ , there exists a sequence  $(x_n)$  in S such that  $x_n \to l$  as  $n \to \infty$ .
  - $\langle 2 \rangle 1$ . Let:  $l \in \operatorname{cl} S$
  - $\langle 2 \rangle 2$ . For  $n \in \mathbb{Z}_+$ , pick  $x_n \in B(l, 1/n) \cap S$

PROOF: There must be such an  $x_n$  otherwise S - B(l, 1/n) would be a smaller closed set that includes S.

 $\langle 2 \rangle 3. \ x_n \to l \text{ as } n \to \infty$ 

 $\langle 1 \rangle 2$ . For any sequence  $(x_n)$  in S, if  $x_n \to l$  as  $n \to \infty$  then  $l \in \operatorname{cl} S$ .

PROOF: Theorem 18.1.40.

**Definition 18.1.44** (Dense). Let V be a normed space over K. Let  $S \subseteq V$ . Then S is dense if and only if cl S = V.

**Theorem 18.1.45** (Weierstrass Approximation Theorem). Let a and b be real numbers with a < b. In C([a,b]), the set of polynomials is dense.

PROOF:TODO

**Proposition 18.1.46.** *Let*  $p \ge 1$ . *The set of all sequences that have only finitely* many non-zero terms is dense in  $l^p$ .

#### Proof:

 $\langle 1 \rangle 1$ . Let:  $(z_n) \in l^p$ 

 $\langle 1 \rangle 2$ . Let:  $\epsilon > 0$ 

PROVE: There exists a sequence  $(x_n)$  with only finitely many non-zero terms such that  $(\sum_{n=1}^{\infty}|z_n-x_n|^p)^{1/p}<\epsilon$   $\langle 1\rangle 3$ . PICK N such that  $|\sum_{n=1}^{\infty}|z_n|^p-\sum_{n=1}^{N}|z_n|^p|<\epsilon^p$   $\langle 1\rangle 4$ . Let:  $(x_n)$  be the sequence that agrees with  $(z_n)$  up to term N, and then

zeros after that.  $\langle 1 \rangle$ 5.  $(\sum_{n=1}^{\infty} |z_n - x_n|^p)^{1/p} < \epsilon$ 

Proof:

$$\left(\sum_{n=1}^{\infty} |z_n - x_n|^p\right)^{1/p} = \left(\sum_{n=N+1}^{\infty} |z_n|^p\right)^{1/p}$$

$$< \epsilon$$

$$(\langle 1 \rangle 4)$$

**Theorem 18.1.47.** Let V be a normed space over K. Let  $S \subseteq V$ . Then the following are equivalent.

- 1. S is dense.
- 2. For all  $l \in V$ , there exists a sequence  $(x_n)$  in S such that  $x_n \to l$  as
- 3. Every nonempty open subset of V intersects S.

#### Proof:

 $\langle 1 \rangle 1$ .  $1 \Leftrightarrow 2$ 

PROOF: Theorem 18.1.43.

- $\langle 1 \rangle 2. \ 1 \Rightarrow 3$ 
  - $\langle 2 \rangle 1$ . Assume: S is dense.
  - $\langle 2 \rangle 2$ . Let: U be a nonempty open subset of V.
  - $\langle 2 \rangle 3$ . X U does not include S.

```
PROOF: Lest we have \operatorname{cl} S \subseteq X - U. \langle 2 \rangle 4. U intersects S. \langle 1 \rangle 3. 3 \Rightarrow 1 \langle 2 \rangle 1. Assume: Every nonempty subset of V intersects S. \langle 2 \rangle 2. Every closed proper subset of V does not include S. \langle 2 \rangle 3. \operatorname{cl} S = V
```

**Definition 18.1.48** (Compact). Let V be a normed space over K and  $S \subseteq V$ . Then S is *compact* if and only if every sequence in S has a convergent subsequence whose limit is in S.

**Proposition 18.1.49.** In  $K^n$ , a set is compact if and only if it is bounded and closed.

PROOF: TODO

**Definition 18.1.50** (Bounded). Let V be a normed space over K and  $S \subseteq V$ . Then S is bounded iff there exists r > 0 such that  $V \subseteq B(0, r)$ .

Theorem 18.1.51. Every compact set is closed and bounded.

```
Proof:
\langle 1 \rangle 1. Let: V be a normed space over K.
\langle 1 \rangle 2. Let: S \subseteq V be compact.
\langle 1 \rangle 3. S is closed.
    \langle 2 \rangle 1. Let: (x_n) be a sequence in S that converges to l
    \langle 2 \rangle 2. PICK a sequence (x_{n_r}) that converges to x \in S
       Proof: \langle 1 \rangle 2, \langle 2 \rangle 1
    \langle 2 \rangle 3. \ x_{n_r} \to l \text{ as } n \to \infty
       Proof: \langle 2 \rangle 1, \langle 2 \rangle 2
    \langle 2 \rangle 4. \ l = x
       Proof: Proposition 18.1.2.
    \langle 2 \rangle 5. \ l \in S
       Proof: \langle 2 \rangle 2, \langle 2 \rangle 4
    \langle 2 \rangle 6. Q.E.D.
       Proof: Theorem 18.1.40.
\langle 1 \rangle 4. S is bounded.
    \langle 2 \rangle 1. Assume: for a contradiction S is unbounded.
    \langle 2 \rangle 2. For n \in \mathbb{Z}_+, PICK x_n \in S - B(0, n)
    \langle 2 \rangle 3. Pick a convergent subsequence (x_{n_r}) that converges to l, say.
    \langle 2 \rangle 4. Pick N \in \mathbb{Z}_+ such that ||l|| < N
    \langle 2 \rangle5. PICK r such that n_r > N and ||x_{n_r} - l|| < N - ||l||
    \langle 2 \rangle 6. \|x_{n_r}\| < N < n_r
    \langle 2 \rangle 7. This contradicts \langle 2 \rangle 2.
```

**Proposition 18.1.52.** In C([0,1]), the closed ball  $\overline{B}(0,1)$  is closed and bounded but not compact.

PROOF: The sequence of functions  $(x^n)$  is in  $\overline{B}(0,1)$  but has no convergent subsequence.  $\square$ 

**Theorem 18.1.53** (Riesz's Lemma). Let V be a normed vector space over K. Let X be a closed proper subspace of V. Let  $0 < \epsilon < 1$ . Then there exists  $x \in V$  such that ||x|| = 1 and  $\forall y \in X. ||x - y|| \ge \epsilon$ .

#### Proof:

$$\langle 1 \rangle 1$$
. Pick  $z \in V - X$ 

$$\langle 1 \rangle 2$$
. Let:  $d = \inf_{x \in X} ||z - x||$ 

$$\langle 1 \rangle 3. \ d > 0$$

PROOF: Since X is closed, there exists e > 0 such that  $B(z,d) \subseteq V - X$  and hence  $||z - x|| \ge d$  for all  $x \in X$ .

 $\langle 1 \rangle 4$ . PICK  $x_0 \in X$  such that  $d \leq ||z - x_0|| \leq d/\epsilon$ 

PROOF: One exists since  $d/\epsilon$  is not a lower bound for  $\{||z-x|| \mid x \in X\}$ .

$$\langle 1 \rangle 5$$
. Let:  $x = (z - x_0) / ||z - x_0||$ 

$$\langle 1 \rangle 6$$
. Let:  $y \in X$ 

$$\langle 1 \rangle 7. ||x - y|| \ge \epsilon$$

Proof:

$$||x - y|| = \left\| \frac{z - x_0}{||z - x_0||} - y \right\|$$

$$= \frac{1}{||z - x_0||} ||z - (x_0 + ||z - x_0||y)||$$

$$\geq \frac{1}{||z - x_0||} d$$

$$\geq \epsilon$$

$$(\langle 1 \rangle 2)$$

$$\geq \epsilon$$

$$(\langle 1 \rangle 4)$$

**Theorem 18.1.54.** Let V be a normed space over K. Then V is finite dimensional if and only if  $\overline{B}(0,1)$  is compact.

- $\langle 1 \rangle 1$ . If V is finite dimensional then  $\overline{B}(0,1)$  is compact.
  - $\langle 2 \rangle 1$ . Assume: V is finite dimensional.
  - $\langle 2 \rangle 2$ . Pick a basis  $\{e_1, \ldots, e_n\}$
  - $\langle 2 \rangle 3$ . Assume: w.l.o.g.  $\|\alpha_1 e_1 + \cdots + \alpha_n e_n\| = |\alpha_1| + \cdots + |\alpha_n|$
  - $\langle 2 \rangle 4$ . Let:  $(\alpha_{k1}e_1 + \cdots + \alpha_{kn}e_n)$  be a sequence in  $\overline{B}(0,1)$
  - $\langle 2 \rangle$ 5. PICK a convergent subsequence  $(\alpha_{k_r 1})$  of  $(\alpha_{k1})$ , a convergent subsequence  $(\alpha_{k'_r} 2)$  of  $(\alpha_{k_r 2}), \ldots,$  a convergent subsequence  $(\alpha_{k''_r} n)$ .
  - $\langle 2 \rangle 6$ .  $(\alpha_{k_r''1}e_1 + \cdots + \alpha_{k_r''n})$  converges.
- $\langle 1 \rangle 2$ . If V is infinite dimensional then  $\overline{B}(0,1)$  is not compact.
  - $\langle 2 \rangle 1$ . Assume: V is infinite dimensional.
  - $\langle 2 \rangle 2$ . Choose a sequence  $(x_n)$  with  $||x_n|| = 1$  and  $||x_m x_n|| \ge 1/2$  for  $m \ne n$ 
    - $\langle 3 \rangle 1$ . Assume:  $x_1, \ldots, x_n$  satisfy  $||x_i|| = 1$  and  $||x_i x_j|| \ge 1/2$  for  $i \ne j$
    - (3)2. PICK  $x_{n+1} \in V$  such that  $||x_{n+1}|| = 1$  and for all  $y \in \text{span}\{x_1, \dots, x_n\}$  we have  $||x_{n+1} y|| \ge 1/2$

```
\langle 4 \rangle 1. span\{x_1, \ldots, x_n\} is closed.
              \langle 5 \rangle 1. Let: S = \operatorname{span}\{x_1, \dots, x_n\}
              \langle 5 \rangle 2. Let: (a_n) be a sequence in S that converges to a \in V
                      Prove: a \in S
              \langle 5 \rangle 3. (a_n) is a Cauchy sequence in V.
              \langle 5 \rangle 4. (a_n) is a Cauchy sequence in S.
              \langle 5 \rangle 5. A finite dimensional normed space is a Banach space.
              \langle 5 \rangle 6. S is complete.
              \langle 5 \rangle 7. \ a \in S
          \langle 4 \rangle 2. span\{x_1, \ldots, x_n\} is a proper subspace of V.
             Proof: \langle 2 \rangle 1
          \langle 4 \rangle3. Q.E.D.
             Proof: Riesz's Lemma.
    \langle 2 \rangle 3. Assume: for a contradiction (x_{n_r}) is a subsequence that converges to l
    \langle 2 \rangle 4. For all r \in \mathbb{N}, we have ||x_{n_r} - l|| + ||x_{n_{r+1}} - l|| \ge 1/2
    \langle 2 \rangle 5. This is a contradiction.
```

**Proposition 18.1.55.** Let V be a normed space. The closure of a subspace of V is a subspace.

```
Proof:
```

```
\langle 1 \rangle1. Let: U be a subspace of V \langle 1 \rangle2. Let: x, y \in \operatorname{cl} U and \alpha, \beta \in K \langle 1 \rangle3. Pick sequences (x_n), (y_n) in U that converge to x and y respectively. \langle 1 \rangle4. \alpha x_n + \beta y_n \to \alpha x + \beta y as n \to \infty \langle 1 \rangle5. \alpha x + \beta y \in \operatorname{cl} U
```

### 18.2 Continuous Linear Mappings

**Definition 18.2.1** (Continuous). Let U and V be normed spaces. Let  $f: U \to V$  and  $x \in U$ . Then f is continuous at x iff, for any sequence  $(x_n)$  in U, if  $x_n \to x$  as  $n \to \infty$  then  $f(x_n) \to f(x)$  as  $n \to \infty$ . The function f is continuous iff f is continuous at every point.

**Proposition 18.2.2.** *Let* V *be a normed space. Then the norm is a continuous function*  $V \to \mathbb{R}$ .

PROOF: From Proposition 18.0.4.

**Proposition 18.2.3.** Let U and V be normed space. Let  $f: U \to V$ . Then the following are equivalent.

- 1. f is continuous.
- 2. For any open set S in V, we have  $f^{-1}(S)$  is open in U.

3. For any closed set C in V, we have  $f^{-1}(C)$  is closed in U.

**Proposition 18.2.4.** Let U and V be normed spaces over K. Let  $T: U \to V$  be a linear transformation. If T is continuous at some point, then it is continuous.

#### Proof:

- $\langle 1 \rangle 1$ . Assume: T is continuous at  $u_0$
- $\langle 1 \rangle 2$ . Let:  $x_n \to l$  as  $n \to \infty$  in U
- $\langle 1 \rangle 3$ .  $x_n l + u_0 \to u_0$  as  $n \to \infty$ .
- $\langle 1 \rangle 4$ .  $T(x_n l + u_0) \to T(u_0)$  as  $n \to \infty$ .
- $\langle 1 \rangle 5$ .  $T(x_n) T(l) + T(u_0) \to T(u_0)$  as  $n \to \infty$ .
- $\langle 1 \rangle 6. \ T(x_n) \to T(l) \text{ as } n \to \infty.$

**Definition 18.2.5** (Bounded). Let U and V be normed spaces over K. Let  $T: U \to V$  be a linear transformation. Then T is bounded iff there exists  $\alpha > 0$  such that, for all  $x \in U$ , we have  $||T(x)|| \le \alpha ||x||$ .

**Theorem 18.2.6.** Let U and V be normed spaces over K. Let  $T: U \to V$  be a linear transformation. Then T is continuous if and only if it is bounded.

#### Proof:

- $\langle 1 \rangle 1$ . If T is continuous then T is bounded.
  - $\langle 2 \rangle 1$ . Assume: T is not bounded.
  - $\langle 2 \rangle 2$ . For  $n \in \mathbb{Z}_+$ , PICK  $x_n \in U$  such that  $||T(x_n)|| > n||x_n||$ .
  - $\langle 2 \rangle 3$ . For  $n \in \mathbb{Z}_+$ , LET:

$$y_n = \frac{x_n}{n||x_n||}$$

- $\langle 2 \rangle 4$ .  $y_n \to 0$  as  $n \to \infty$
- $\langle 2 \rangle 5$ .  $T(y_n) \not\to 0$  as  $n \to \infty$
- $\langle 2 \rangle 6$ . T is not continuous.
- $\langle 1 \rangle 2$ . If T is bounded then T is continuous.
  - $\langle 2 \rangle 1$ . Assume: T is bounded.
  - $\langle 2 \rangle 2$ . PICK  $\alpha > 0$  such that  $\forall x \in U ||T(x)|| \leq \alpha ||x||$ .
  - $\langle 2 \rangle 3$ . T is continuous at 0.
    - $\langle 3 \rangle 1$ . Let:  $(x_n)$  be a sequence that converges to 0 in U
    - $\langle 3 \rangle 2$ .  $T(x_n) \to 0$  as  $n \to \infty$

Proof:

$$||T(x_n)|| \le \alpha ||x_n||$$
  $(\langle 2 \rangle 2)$   
  $\to 0$  as  $n \to \infty$ 

 $\langle 2 \rangle 4$ . T is continuous.

Proof: Proposition 18.2.4.

**Proposition 18.2.7.** Let U and V be normed spaces over K where U is finite dimensional. Let  $T: U \to V$  be a linear transformation. Then T is bounded.

Proof:

- $\langle 1 \rangle 1$ . PICK a basis  $\{e_1, \ldots, e_n\}$  of unit vectors for U.
- $\langle 1 \rangle 2$ . Let:  $M = \max(||T(e_1)||, \dots, ||T(e_n)||)$
- $\langle 1 \rangle 3$ . Pick C > 0 such that, for all  $\alpha_1, \ldots, \alpha_n \in K$ , we have  $|\alpha_1| + \cdots + |\alpha_n| \leq$  $C\|\alpha_1e_1+\cdots+\alpha_ne_n\|$

PROOF: Theorem 18.1.13.

 $\langle 1 \rangle 4$ . Let:  $x \in U$ 

PROVE:  $||T(x)|| \le CM||x||$ 

- $\langle 1 \rangle 5$ . Let:  $x = \alpha_1 e_1 + \cdots + \alpha_n e_n$
- $\langle 1 \rangle 6$ .  $||T(x)|| \leq CM||x||$

Proof:

$$||T(x)|| = ||\alpha_1 T(e_1) + \dots + \alpha_n T(e_n)||$$
 (T linear)  

$$\leq |\alpha_1| ||T(e_1)|| + \dots + |\alpha_n| ||T(e_n)||$$
 (Triangle inequality)  

$$\leq M(|\alpha_1| + \dots + |\alpha_n|)$$
 (\langle 1\lambda)  

$$\leq CM||x||$$
 (\lambda 1\lambda)

 $(\langle 1 \rangle 3)$ 

Corollary 18.2.7.1. Let U and V be normed spaces over K where U is finite dimensional. Let  $T: U \to V$  be a linear transformation. Then T is continuous.

**Proposition 18.2.8.** Let U and V be normed spaces over K. Let  $T: U \to V$ be a linear transformation. If T is continuous, then T is uniformly continuous.

#### Proof:

- $\langle 1 \rangle 1$ . Assume: T is continuous
- $\langle 1 \rangle 2$ . Pick B > 0 such that  $\forall x \in U ||T(x)|| \leq B||x||$
- $\langle 1 \rangle 3$ . Let:  $\epsilon > 0$
- $\langle 1 \rangle 4$ . Let:  $\delta = \epsilon/B$
- $\langle 1 \rangle 5$ . Let:  $x, y \in U$
- $\langle 1 \rangle 6$ . Assume:  $||x y|| < \delta$
- $\langle 1 \rangle 7$ .  $||T(x) T(y)|| < \epsilon$

Proof:

$$||T(x) - T(y)|| = ||T(x - y)||$$

$$\leq B||x - y||$$

$$< B\delta$$

$$= \epsilon$$

$$(\langle 1 \rangle 2)$$

$$(\langle 1 \rangle 6)$$

$$(\langle 1 \rangle 4)$$

**Proposition 18.2.9.** Let U and V be normed spaces over K. The set  $\mathcal{B}(U,V)$ of all bounded linear maps from U to V forms a subspace of the space of all linear maps from U to V.

- $\langle 1 \rangle 1$ . Let:  $S, T : U \to V$  be bounded linear maps and  $\alpha, \beta \in K$ . PROVE:  $\alpha S + \beta T$  is bounded.
- $\langle 1 \rangle 2$ . PICK B, C > 0 such that  $\forall x \in U ||S(x)|| \leq B||x||$  and  $||T(x)|| \leq C||x||$
- $\langle 1 \rangle 3. \ \forall x \in U. \|(\alpha S + \beta T)(x)\| \le (|\alpha|B + |\beta|C)\|x\|$

**Proposition 18.2.10.** Let U and V be normed spaces over K. Define the operator norm  $\| \|$  on  $\mathcal{B}(U,V)$  by  $\|T\| := \sup\{\|T(x)\| \mid x \in U, \|x\| = 1\}$ . Then  $\| \|$  is a norm on  $\mathcal{B}(U,V)$ .

#### Proof:

```
\langle 1 \rangle 1. For all T \in \mathcal{B}(U, V), the set \{ ||T(x)|| \mid x \in U, ||x|| = 1 \} is bounded above.
```

$$\langle 2 \rangle 1$$
. Let:  $T \in \mathcal{B}(U, V)$ 

$$\langle 2 \rangle 2$$
. Pick B such that  $\forall x \in U . ||T(x)|| \leq B||x||$ .

$$\langle 2 \rangle 3$$
. *B* is an upper bound for  $\{ ||T(x)|| \mid x \in U, ||x|| = 1 \}$ .

$$\langle 1 \rangle 2$$
. If  $||T|| = 0$  then  $T = 0$ .

$$\langle 2 \rangle 1$$
. Assume:  $||T|| = 0$ 

$$\langle 2 \rangle 2$$
. Let:  $x \in U$ 

PROVE: 
$$T(x) = 0$$

$$\langle 2 \rangle 3$$
. Assume: w.l.o.g.  $||x|| \neq 0$ 

$$\langle 2 \rangle 4$$
. Let:  $y = x/||x||$ 

$$\langle 2 \rangle 5$$
.  $||y|| = 1$ 

$$\langle 2 \rangle 6. \ \|T(y)\| = 0$$

$$\langle 2 \rangle 7$$
.  $T(y) = 0$ 

$$\langle 2 \rangle 8. \ T(x) = 0$$

$$\langle 1 \rangle 3$$
. For all  $\lambda \in K$  and  $T \in \mathcal{B}(U,V)$ , we have  $\|\lambda T\| = |\lambda| \|T\|$ 

$$\langle 2 \rangle 1$$
. Let:  $\lambda \in K$  and  $T \in \mathcal{B}(U, V)$ 

$$\langle 2 \rangle 2$$
.  $||\lambda T|| = |\lambda|||T||$ 

Proof:

$$\begin{split} \|\lambda T\| &= \sup\{\|\lambda T(x)\| \mid x \in U, \|x\| = 1\} \\ &= \sup\{|\lambda| \|T(x)\| \mid x \in U, \|x\| = 1\} \\ &= |\lambda| \sup\{\|T(x)\| \mid x \in U, \|x\| = 1\} \\ &= |\lambda| \|T\| \end{split}$$

 $\langle 1 \rangle 4$ . For all  $S, T \in \mathcal{B}(U, V)$ , we have  $||S + T|| \le ||S|| + ||T||$ .

$$\langle 2 \rangle 1$$
. Let:  $S, T \in \mathcal{B}(U, V)$ 

$$\langle 2 \rangle 2. \|S + T\| \le \|S\| + \|T\|$$

Proof:

$$\begin{split} \|S+T\| &= \sup\{\|S(x)+T(x)\| \mid x \in U, \|x\| = 1\} \\ &\leq \sup\{\|S(x)\| + \|T(x)\| \mid x \in U, \|x\| = 1\} \\ &\leq \sup\{\|S(x)\| \mid x \in U, \|x\| = 1\} + \sup\{\|T(x)\| \mid x \in U, \|x\| = 1\} \\ &= \|S\| + \|T\| \end{split}$$

**Proposition 18.2.11.** Let U and V be normed spaces. Let  $T \in \mathcal{B}(U,V)$ . Then ||T|| is the least number such that  $\forall u \in U.||T(u)|| \leq ||T|| ||u||$ .

$$\langle 1 \rangle 1. \ \forall u \in U. ||T(u)|| \le ||T|| ||u||$$

$$\langle 2 \rangle 1$$
. Let:  $u \in U$ 

$$\langle 2 \rangle 2$$
. Let:  $v = u/\|u\|$ 

```
 \begin{array}{l} \langle 2 \rangle 3. \ \|T(v)\| \leq \|T\| \\ \langle 2 \rangle 4. \ \|T(u)\| \leq \|T\| \|u\| \\ \langle 1 \rangle 2. \ \text{If } \alpha \ \text{satisfies} \ \forall u \in U. \|T(u)\| \leq \alpha \|u\|, \ \text{then} \ \|T\| \leq \alpha \\ \langle 2 \rangle 1. \ \text{Assume:} \ \forall u \in U. \|T(u)\| \leq \alpha \|u\| \\ \langle 2 \rangle 2. \ \text{For all} \ x \in U, \ \text{if} \ \|x\| = 1 \ \text{then} \ \|T(x)\| \leq \alpha \\ \langle 2 \rangle 3. \ \|T\| \leq \alpha \\ \hline \\ \\ \\ \\ \\ \\ \\ \\ \end{array}
```

**Proposition 18.2.12.** Let V be a normed space. Then  $id_V$  is a bounded linear function  $V \to V$ , and  $\|id_V\| = 1$ .

**Proposition 18.2.13.** Let U and V be normed spaces. The constant zero function  $U \to V$  is a bounded linear transformation with norm 0.

**Proposition 18.2.14.** Let  $N \in \mathbb{N}$ . Let  $T : \mathbb{C}^N \to \mathbb{C}^N$  be a linear transformation with matrix  $A = (a_{ij})$ . Then T is bounded and

$$||T|| \le \sqrt{\sum_{i=1}^{N} \sum_{j=1}^{N} |a_{ij}|^2}$$
.

**Definition 18.2.15** (Uniform Convergence). Let U and V be normed spaces. Let  $(T_n)$  be a sequence in  $\mathcal{B}(U,V)$  and  $T \in \mathcal{B}(U,V)$ . Then  $(T_n)$  converges uniformly to T iff  $(T_n)$  converges to T under the standard norm defined above.

**Theorem 18.2.16.** Let U and V be normed spaces. Let  $T: U \to V$  be a continuous linear function. Then  $\ker T$  is a closed subspace of U.

#### Proof:

 $\langle 1 \rangle 1$ . ker T is a subspace of U

PROOF: If  $x, y \in \ker T$  and  $\alpha, \beta \in K$  then  $T(\alpha x + \beta y) = \alpha T(x) + \beta T(y) = 0$ .  $\langle 1 \rangle 2$ .  $\ker T$  is closed.

PROOF: Let  $(x_n)$  be a sequence in ker T and  $x_n \to l$ . Then  $T(l) = \lim_{n \to \infty} T(x_n) = 0$ .

**Theorem 18.2.17.** Let U and V be normed spaces. Let W be a closed subspace of U and  $T: W \to V$  be a continuous linear mapping. Then the graph  $G = \{(x, T(x)) \mid x \in W\}$  is closed in  $U \times V$ .

- $\langle 1 \rangle 1$ . Assume: w.l.o.g.  $T \neq 0$
- $\langle 1 \rangle 2$ . Let:  $(x,y) \in (U \times V) G$ , i.e.  $y \neq T(x)$
- $\langle 1 \rangle 3$ . Let:  $\epsilon = ||y T(x)|| > 0$
- $\langle 1 \rangle 4$ . Let:  $x' \in W$  with  $||x x'|| < \epsilon/3||T||$
- $\langle 1 \rangle 5$ . Let:  $y' \in V$  with  $||y y'|| < \epsilon/3$
- $\langle 1 \rangle 6. \ y' \neq T(x')$

Proof:

$$||y' - T(x')|| \ge ||y - T(x)|| - ||y - y'|| - ||T(x) - T(x')||$$

$$> \epsilon/3$$

$$> 0$$

**Theorem 18.2.18** (Diagonal Theorem). Let E be a normed space over K. Let  $(x_{ij})$  be an infinite matrix of elements of V. If:

- 1. For all  $j \in \mathbb{Z}_+$ , we have  $x_{ij} \to 0$  as  $i \to \infty$ ;
- 2. Every increasing sequence of positive integers  $(p_j)$  has a subsequence  $(p_{j_r})$ such that

$$\sum_{s=1}^{\infty} x_{p_{j_r}p_{j_s}} \to 0 \text{ as } r \to \infty$$

then  $x_{ii} \to 0$  as  $i \to \infty$ .

- $\langle 1 \rangle 1$ . Assume: for a contradiction  $x_{ii} \not\to 0$  as  $i \to \infty$
- $\langle 1 \rangle 2$ . PICK  $\epsilon > 0$  such that, for all N, there exists  $n \geq N$  such that  $||x_{nn}|| \geq \epsilon$
- $\langle 1 \rangle 3$ . PICK an increasing sequence of integers  $(p_j)$  such that  $||x_{p_jp_j}|| \geq \epsilon$  for all j.
- $\langle 1 \rangle 4$ . PICK a subsequence  $(q_i)$  such that  $\sum_{j=1}^{\infty} x_{q_i q_j} \to 0$  as  $i \to \infty$
- $\langle 1 \rangle$ 5. For all i, we have  $x_{q_i q_j} \to 0$  as  $j \to \infty$   $\langle 1 \rangle$ 6. For all j, we have  $x_{q_i q_j} \to 0$  as  $i \to \infty$
- $\langle 1 \rangle 7$ . Define a subsequence  $(r_n)$  of  $(q_i)$  by  $r_1 = q_1$  and, for all  $n, r_{n+1}$  is the first entry such that  $r_{n+1} > r_n$ ,  $||x_{q_i r_n}|| < \epsilon/2^{n+1}$  for all  $q_i \ge r_{n+1}$ , and  $||x_{r_j r_{n+1}}|| < \epsilon/2^{n+2}$  for  $j = 1, \ldots, n$ .
- $\langle 1 \rangle 8$ .  $||x_{r_i r_j}|| < \epsilon/2^{j+1}$  for all i, j such that  $i \neq j$   $\langle 1 \rangle 9$ . PICK a subsequence  $(s_j)$  of  $(r_j)$  such that  $\sum_{j=1}^{\infty} x_{s_i s_j} \to 0$  as  $i \to \infty$   $\langle 1 \rangle 10$ . For all i we have  $||\sum_{j=1}^{\infty} x_{s_i s_j}|| \geq \epsilon/2$

Proof

$$\left\| \sum_{j=1}^{\infty} x_{s_i s_j} \right\| = \left\| x_{s_i s_i} + \sum_{i \neq j} x_{s_i s_j} \right\|$$

$$\geq \left\| \|x_{s_i s_i}\| - \left\| \sum_{i \neq j} x_{s_i s_j} \right\|$$

$$\geq \left\| \|x_{s_i s_i}\| - \sum_{i \neq j} \|x_{s_i s_j}\| \right\|$$

$$\geq \epsilon/2$$

$$(\langle 1 \rangle 2, \langle 1 \rangle 8)$$

 $\langle 1 \rangle 11$ . Q.E.D.

PROOF:  $\langle 1 \rangle 9$  and  $\langle 1 \rangle 10$  form a contradiction.

### 18.3 Banach Spaces

**Definition 18.3.1** (Cauchy Sequence). Let V be a normed space over K. A Cauchy sequence is a sequence of points  $(x_n)$  such that, for every  $\epsilon > 0$ , there exists N such that  $\forall m, n \geq N . ||x_m - x_n|| < \epsilon$ .

**Theorem 18.3.2.** Let V be a normed space over K. Let  $(x_n)$  be a sequence in V. The following are equivalent.

- 1.  $(x_n)$  is Cauchy.
- 2. For every pair of increasing sequences of positive integers  $(p_n)$  and  $(q_n)$ , we have  $||x_{p_n} x_{q_n}|| \to 0$  as  $n \to \infty$ .
- 3. For every increasing sequence of positive integers  $(p_n)$ , we have  $||x_{p_n} x_n|| \to 0$  as  $n \to \infty$ .

#### Proof:

- $\langle 1 \rangle 1. \ 1 \Rightarrow 2$ 
  - $\langle 2 \rangle 1$ . Assume:  $(x_n)$  is Cauchy.
  - $\langle 2 \rangle 2$ . Let:  $(p_n)$  and  $(q_n)$  are increasing sequences of positive integers.
  - $\langle 2 \rangle 3$ . Let:  $\epsilon > 0$
  - $\langle 2 \rangle 4$ . PICK N such that  $\forall m, n \geq N . ||x_m x_n|| < \epsilon$
  - $\langle 2 \rangle$ 5.  $\forall n \geq N. ||x_{p_n} x_{q_n}|| < \epsilon$ PROOF: Since  $p_n, q_n \geq n \geq N$ .
- $\langle 1 \rangle 2. \ 2 \Rightarrow 3$

PROOF: Trivial.

- $\langle 1 \rangle 3. \ 2 \Rightarrow 1$ 
  - $\langle 2 \rangle 1$ . Assume:  $(x_n)$  is not Cauchy
  - $\langle 2 \rangle 2$ . Pick  $\epsilon > 0$  such that, for every  $N \in \mathbb{Z}_+$ , there exist  $m_N, n_N \geq N$  such that  $||x_{m_N} x_{n_N}|| \geq \epsilon$
  - $\langle 2 \rangle 3$ . Assume: w.l.o.g.  $(m_N)$  and  $(n_N)$  are increasing sequences.
  - $\langle 2 \rangle 4$ .  $||x_{m_N} x_{n_N}|| \not\to 0$  as  $N \to \infty$ .
- $\langle 1 \rangle 4. \ 3 \Rightarrow 2$ 
  - $\langle 2 \rangle 1$ . Assume: 3
  - $\langle 2 \rangle 2$ . Let:  $(p_n)$  and  $(q_n)$  be increasing sequences of positive integers.
  - $\langle 2 \rangle 3$ . Let:  $\epsilon > 0$
  - $\langle 2\rangle 4.$  Pick N such that  $\forall n\geq N.\|x_{p_n}-x_n\|<\epsilon/2$  and  $\forall n\geq N.\|x_{q_n}-x_n\|<\epsilon/2$
- $\langle 2 \rangle 5. \ \forall n \ge N. \|x_{p_n} x_{q_n}\| < \epsilon$

Proposition 18.3.3. Every convergent sequence is Cauchy.

- $\langle 1 \rangle 1$ . Let:  $x_n \to l$  as  $n \to \infty$ .
- $\langle 1 \rangle 2$ . Let:  $\epsilon > 0$
- $\langle 1 \rangle 3$ . PICK N such that  $\forall n \geq N . ||x_n l|| < \epsilon/2$

 $\langle 1 \rangle 4$ . For all  $m, n \geq N$  we have  $||x_m - x_n|| < \epsilon$ .

Proposition 18.3.4. In  $\mathcal{P}([0,1])$ , let

$$P_n(x) = 1 + x + \frac{x^2}{2} + \dots + \frac{x^n}{n!}$$
.

Then  $(P_n)$  is Cauchy but does not converge.

PROOF: It converges to  $e^x$  in  $\mathcal{C}([0,1])$ , therefore it is Cauchy in  $\mathcal{C}([0,1])$ , hence Cauchy in  $\mathcal{P}([0,1])$ . Since  $e^x \notin \mathcal{P}([0,1])$ , it does not converge in that space.  $\sqcup$ 

**Proposition 18.3.5.** Let V be a normed space over K. Let  $(x_n)$  be a Cauchy sequence in V. Then  $(\|x_n\|)$  converges in  $\mathbb{R}$ .

Proof:

- $\langle 1 \rangle 1$ . ( $||x_n||$ ) is Cauchy.
  - $\langle 2 \rangle 1$ . Let:  $\epsilon > 0$
  - $\langle 2 \rangle 2$ . PICK N such that  $\forall m, n \geq N . ||x_m x_n|| < \epsilon$
  - $\langle 2 \rangle 3. \ \forall m, n \geq N. ||x_m|| ||x_n||| < \epsilon$

Proof: Proposition 18.0.4.

 $\langle 1 \rangle 2$ . Q.E.D.

PROOF: Since every Cauchy sequence in  $\mathbb{R}$  converges.

**Proposition 18.3.6.** Every Cauchy sequence is bounded.

Proof:

- $\langle 1 \rangle 1$ . Let: V be a normed space over K.
- $\langle 1 \rangle 2$ . Let:  $(x_n)$  be a Cauchy sequence in V.
- $\langle 1 \rangle 3$ . PICK N such that  $\forall m, n \geq N . ||x_m x_n|| < 1$ .
- $\langle 1 \rangle 4$ . Let:  $B = \max(||x_1||, \dots, ||x_{N-1}||, ||x_N|| + 1)$
- $\langle 1 \rangle 5. \ \forall n. ||x_n|| \le B$

a Banach space iff every Cauchy sequence converges.

Proposition 18.3.8.  $l^2$  is complete.

Proof:

- $\langle 1 \rangle 1$ . Let:  $(a_n)$  be a Cauchy sequence in  $l^2$  where  $a_n = (\alpha_{n1}, \alpha_{n2}, \ldots)$ .  $\langle 1 \rangle 2$ . For all  $\epsilon > 0$ , there exists  $n_0 \in \mathbb{Z}_+$  such that  $\forall m, n \geq n_0$ .  $\sum_{k=1}^{\infty} |\alpha_{mk} \alpha_{mk}| = 1$

**Definition 18.3.7** (Banach Space). A normed space V over K is complete or

- $\langle 1 \rangle 3$ . For every  $k \in \mathbb{Z}_+$  and  $\epsilon > 0$ , there exists  $n_0 \in \mathbb{Z}_+$  such that  $\forall m, n \geq 1$  $n_0.|\alpha_{mk}-\alpha_{nk}|<\epsilon.$
- $\langle 1 \rangle 4$ . For every  $k \in \mathbb{Z}_+$ ,  $(\alpha_{nk})$  is Cauchy in  $\mathbb{C}$ .
- $\langle 1 \rangle 5$ . For every  $k \in \mathbb{Z}_+$ ,  $(\alpha_{nk})$  converges in  $\mathbb{C}$ .
- $\langle 1 \rangle 6$ . For  $k \in \mathbb{Z}_+$ ,

```
Let: \alpha_k = \lim_{n \to \infty} \alpha_{nk}
\langle 1 \rangle 7. Let a be the sequence (\alpha_k)
(1)8. For all \epsilon > 0, there exists n_0 such that \forall n \geq n_0. \sum_{k=1}^{\infty} |\alpha_k - \alpha_{nk}|^2 \leq \epsilon^2.
   PROOF: Letting m \to \infty in \langle 1 \rangle 2.
\langle 1 \rangle 9. \ a \in l^2
    \langle 2 \rangle 1. PICK n_0 such that \forall n \geq n_0. \sum_{k=1}^{\infty} |\alpha_k - \alpha_{nk}|^2 \leq 1
    \langle 2 \rangle 2. \ (\alpha_k - \alpha_{n_0 k}) \in l^2
\langle 2 \rangle 3. \ (\alpha_{n_0 k}) \in l^2
       PROOF: By \langle 1 \rangle 1 since the sequence is a_{n_0}.
    \langle 2 \rangle 4. \ (\alpha_k) \in l^2
       Proof: Proposition 16.0.2.
\langle 1 \rangle 10. \ a_n \to a \text{ as } n \to \infty
   PROOF: By \langle 1 \rangle 8 since ||a - a_n|| = \sqrt{\sum_{k=1}^{\infty} |\alpha_k - \alpha_{nk}|^2}.
Proposition 18.3.9. Let a and b be real numbers with a < b. The space C([a,b])
is complete.
Proof:
\langle 1 \rangle 1. Let: X = [a, b]
\langle 1 \rangle 2. Let: (f_n) be a Cauchy sequence in \mathcal{C}([a,b]).
\langle 1 \rangle 3. For all \epsilon > 0, there exists n_0 such that \forall n, m \geq n_0 . ||f_n - f_m|| < \epsilon.
\langle 1 \rangle 4. For all \epsilon > 0, there exists n_0 such that \forall n, m \geq n_0 . \forall x \in X. | f_n(x) - f_n(x)| = 0
          |f_m(x)| < \epsilon.
\langle 1 \rangle 5. For all x \in [a, b], (f_n(x)) is Cauchy.
\langle 1 \rangle 6. Define f: [a,b] \to \mathbb{C} by f(x) = \lim_{n \to \infty} f_n(x).
\langle 1 \rangle 7. For all \epsilon > 0, there exists n_0 such that \forall n \geq n_0 . \forall x \in X . |f_n(x) - f(x)| < \epsilon
   PROOF: Letting m \to \infty in \langle 1 \rangle 4.
\langle 1 \rangle 8. f is continuous
    \langle 2 \rangle 1. Let: x_0 \in X
    \langle 2 \rangle 2. Let: \epsilon > 0
    \langle 2 \rangle 3. PICK n_0 such that \forall n \geq n_0 . \forall x \in X . |f_n(x) - f(x)| < \epsilon/3
       PROOF: By \langle 1 \rangle 7.
    \langle 2 \rangle 4. Pick \delta > 0 such that \forall x \in X | |x - x_0| < \delta \Rightarrow |f_{n_0}(x) - f_{n_0}(x_0)| < \epsilon/3
       PROOF: Since f_{n_0} is continuous.
    \langle 2 \rangle 5. Let: x \in X
    \langle 2 \rangle 6. Assume: |x - x_0| < \delta
    \langle 2 \rangle 7. |f(x) - f(x_0)| < \epsilon
       Proof:
       |f(x) - f(x_0)| \le |f(x) - f_{n_0}(x)| + |f_{n_0}(x) - f_{n_0}(x_0)| + |f_{n_0}(x_0) - f(x_0)| (Triangle Inequality)
                                 <\epsilon/3+\epsilon/3+\epsilon/3
                                                                                                                                                       (\langle 2 \rangle 3, \langle 2 \rangle 4)
\langle 1 \rangle 9. (f_n) converges to f uniformly.
    Proof: From \langle 1 \rangle 7
```

**Definition 18.3.10** (Series). Let V be a normed space over K. A convergent series in V is a sequence  $(x_n)$  in V such that  $(x_1 + \cdots + x_n)$  is a convergent sequence, in which case we write  $\sum_{n=1}^{\infty} x_n$  for its limit.

**Definition 18.3.11** (Absolutely Convergent Series). Let V be a normed space over K. An absolutely convergent series in V is a sequence  $(x_n)$  such that  $\sum_{n=1}^{\infty} ||x_n|| < \infty.$ 

**Proposition 18.3.12.** In  $\mathcal{P}([0,1])$ , the series  $\sum_{n=0}^{\infty} x^n/n!$  is absolutely convergent but not convergent.

Proof: Proposition 18.3.4.

**Theorem 18.3.13.** A normed space is complete if and only if every absolutely convergent series is convergent.

#### Proof:

- $\langle 1 \rangle 1$ . Let: V be a normed space over K.
- $\langle 1 \rangle 2$ . If V is complete then every absolutely convergent series is convergent.

  - $\langle 2 \rangle 1$ . Assume: V is complete.  $\langle 2 \rangle 2$ . Let:  $\sum_{n=1}^{\infty} a_n$  be absolutely convergent in V.  $\langle 2 \rangle 3$ . For  $n \in \mathbb{Z}_+$ , Let:  $s_n = \sum_{k=1}^n a_k$
  - $\langle 2 \rangle 4$ .  $(s_n)$  is Cauchy.
    - $\langle 3 \rangle 1$ . Let:  $\epsilon > 0$
    - $\langle 3 \rangle$ 2. PICK k such that  $\sum_{n=k+1}^{\infty} ||a_n|| < \epsilon$
    - $\langle 3 \rangle 3$ . Let: m > n > k
    - $\langle 3 \rangle 4$ .  $||s_m s_n|| < \epsilon$

$$||s_m - s_n|| = \left\| \sum_{i=n+1}^m a_i \right\|$$

$$\leq \sum_{i=s+1}^m ||a_i||$$
(Triangle inequality)
$$\leq \sum_{i=k+1}^\infty ||a_i||$$

$$< \epsilon$$
(\lambda 3\rangle 2, \lambda 3\rangle 3)

- $\langle 2 \rangle 5$ .  $(s_n)$  converges.
- $\langle 1 \rangle 3$ . If every absolutely convergent series is convergent then V is complete.
  - $\langle 2 \rangle 1$ . Assume: Every absolutely convergent series in V is convergent.
  - $\langle 2 \rangle 2$ . Let:  $(a_n)$  be a Cauchy sequence in V.
  - $\langle 2 \rangle 3$ . PICK a strictly increasing sequence of positive integers  $(p_n)$  such that  $\forall k. \forall m, n \ge p_k. ||x_m - x_n|| < 2^{-k}.$
  - $\langle 2 \rangle 4$ .  $\sum_{k=1}^{\infty} (x_{p_{k+1}} x_{p_k})$  is absolutely convergent.

$$\sum_{k=1}^{\infty} \|x_{p_{k+1}} - x_{p_k}\| < \sum_{k=1}^{\infty} 2^{-k}$$
 (\langle 2\rangle 3)

$$\langle 2 \rangle 5$$
.  $\sum_{k=1}^{\infty} (x_{p_{k+1}} - x_{p_k})$  is convergent. PROOF:  $\langle 2 \rangle 1$ ,  $\langle 2 \rangle 5$ 

$$\langle 2 \rangle 6$$
. Let:  $s = \sum_{k=1}^{\infty} (x_{p_{k+1}} - x_{p_k})$ 

$$\langle 2 \rangle 7$$
.  $x_{p_k} \to s + x_{p_1}$  as  $k \to \infty$ .

PROOF: 
$$\langle 2 \rangle 1$$
,  $\langle 2 \rangle 3$   
 $\langle 2 \rangle 6$ . Let:  $s = \sum_{k=1}^{\infty} (x_{p_{k+1}} - x_{p_k})$   
 $\langle 2 \rangle 7$ .  $x_{p_k} \to s + x_{p_1}$  as  $k \to \infty$ .  
 $\langle 3 \rangle 1$ .  $\sum_{i=1}^{k} (x_{p_{i+1}} - x_{p_i}) \to s$  as  $k \to \infty$   
PROOF:  $\langle 2 \rangle 6$ 

$$\langle 3 \rangle 2$$
.  $x_{p_{k+1}} - x_{p_1} \to s \text{ as } k \to \infty$ 

$$\langle 2 \rangle 8. \ x_n \to s + x_{p_1} \text{ as } k \to \infty.$$

Proof:

 $\langle 3 \rangle 1$ . Let:  $\epsilon > 0$ 

 $\langle 3 \rangle 2$ . PICK N such that  $\forall k \geq N . ||x_{p_k} - (s + x_{p_1})|| < \epsilon/2$  and  $\forall m, n \geq 1$  $N.||x_m - x_n|| < \epsilon/2$ 

Proof:  $\langle 2 \rangle 2$ ,  $\langle 2 \rangle 7$ 

 $\langle 3 \rangle 3. \ \forall n \geq N. \|x_n - (s + x_{p_1})\| < \epsilon$ 

Theorem 18.3.14. A closed vector subspace of a Banach space is a Banach space.

#### Proof:

- $\langle 1 \rangle 1$ . Let: V be a Banach space over K.
- $\langle 1 \rangle 2$ . Let: U be a closed vector subspace of V.
- $\langle 1 \rangle 3$ . Let:  $(a_n)$  be a Cauchy sequence in U.
- $\langle 1 \rangle 4$ .  $(a_n)$  is a Cauchy sequence in V.
- $\langle 1 \rangle 5$ . Let:  $l = \lim_{n \to \infty} a_n$
- $\langle 1 \rangle 6. \ l \in U$

Proof: Theorem 18.1.40.

 $\langle 1 \rangle 7$ .  $a_n \to l$  as  $n \to \infty$  in U.

**Definition 18.3.15** (Completion). Let V be a normed space over K. A completion of V consists of a normed space W over K and an injection  $\phi: V \to W$ such that:

- 1.  $\forall x, y \in V. \forall \alpha, \beta \in K. \phi(\alpha x + \beta y) = \alpha \phi(x) + \beta \phi(y)$
- 2.  $\forall x \in V || \phi(x) || = ||x||$
- 3.  $\phi(V)$  is dense in W.
- 4. W is complete.

**Proposition 18.3.16.** Every normed space has a completion.

- $\langle 1 \rangle 1$ . Let: V be a normed space over K.
- $\langle 1 \rangle 2$ . Let us say two Cauchy sequences  $(x_n)$ ,  $(y_n)$  ore equivalent iff  $x_n y_n \to 0$  as  $n \to \infty$ .
- $\langle 1 \rangle 3$ . Equivalence is an equivalence relation on the set of Cauchy sequences.
- $\langle 1 \rangle 4$ . Let: W be the set of Cauchy sequences in V quotiented by equivalence.
- $\langle 1 \rangle$ 5. Define addition and multiplication on W by

$$[(x_n)] + [(y_n)] = [(x_n + y_n)]$$
  
 $\lambda[(x_n)] = [(\lambda x_n)]$ 

- $\langle 1 \rangle 6$ . Define a norm on W by  $||[(x_n)]|| = \lim_{n \to \infty} ||x_n||$
- $\langle 1 \rangle 7$ . Define  $\phi: V \to W$  by  $\phi(v) = [(v)]$ .
- $\langle 1 \rangle 8$ .  $\phi$  is injective.
- $\langle 1 \rangle 9. \ \forall x, y \in V. \forall \alpha, \beta \in K. \phi(\alpha x + \beta y) = \alpha \phi(x) + \beta \phi(y)$
- $\langle 1 \rangle 10. \ \forall x \in V. \| \phi(x) \| = \| x \|$
- $\langle 1 \rangle 11$ .  $\phi(V)$  is dense in W.
  - $\langle 2 \rangle 1$ . Let:  $[(a_n)] \in W$  and  $\epsilon > 0$ .

PROVE:  $B([(a_n)], \epsilon)$  intersects  $\phi(V)$ .

- $\langle 2 \rangle 2$ . PICK N such that  $\forall m, n \geq N . ||a_m a_n|| < \epsilon/2$
- $\langle 2 \rangle 3. \ \phi(a_N) \in B([(a_n)], \epsilon)$

Proof:

$$\|[(a_n)] - \phi(a_N)\| = \lim_{n \to \infty} \|a_n - a_N\|$$

$$\leq \epsilon/2$$

$$< \epsilon$$

$$(\langle 2 \rangle 2)$$

- $\langle 1 \rangle 12$ . W is complete.
  - $\langle 2 \rangle 1$ . Let:  $(X_n)$  be a Cauchy sequence in W.
  - $\langle 2 \rangle 2$ . For  $n \in \mathbb{Z}_+$ , PICK  $x_n \in V$  such that

$$\|\phi(x_n) - X_n\| < 1/n$$
.

- $\langle 2 \rangle 3$ .  $(x_n)$  is Cauchy in V.
  - $\langle 3 \rangle 1$ . Let:  $\epsilon > 0$
  - $\langle 3 \rangle 2$ . PICK N such that  $\forall m, n \geq N . ||X_n X_m|| < \epsilon/3$  and  $1/N < \epsilon/3$
  - $\langle 3 \rangle 3$ . Let:  $m, n \geq N$
  - $\langle 3 \rangle 4$ .  $||x_m x_n|| < \epsilon$

Proof:

$$||x_m - x_n|| = ||\phi(x_m) - \phi(x_n)||$$

$$\leq ||\phi(x_m) - X_m|| + ||X_m - X_n|| + ||X_n - \phi(x_n)||$$

$$\leq ||X_m - X_n|| + 1/m + 1/n$$

$$\leq \epsilon$$

- $\langle 2 \rangle 4$ . Let:  $X = [(x_n)]$
- $\langle 2 \rangle 5. \ X_n \to X \text{ as } n \to \infty$

$$||X_n - X|| \le ||X_n - \phi(x_n)|| + ||\phi(x_n) - X||$$
  
 $< ||\phi(x_n) - X|| + 1/n$   
 $\to 0$  as  $n \to \infty$ 

**Proposition 18.3.17.** Let U be a normed space and V a Banach space. Then  $\mathcal{B}(U,V)$  is a Banach space.

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $(T_n)$  be a Cauchy sequence in  $\mathcal{B}(U,V)$
- $\langle 1 \rangle 2$ . For all  $u \in U$ ,  $(T_n(u))$  is a Cauchy sequence in V.
  - $\langle 2 \rangle 1$ . Let:  $u \in U$
  - $\langle 2 \rangle 2$ . Let:  $\epsilon > 0$

PROVE: 
$$\exists N. \forall m, n \geq N. ||T_m(u) - T_n(u)|| < \epsilon$$

- $\langle 2 \rangle 3$ . Assume: w.l.o.g.  $u \neq 0$
- $\langle 2 \rangle 4$ . PICK N such that  $\forall m, n \geq N . ||T_m T_n|| < \epsilon / ||u||$
- $\langle 2 \rangle 5$ . Let:  $m, n \geq N$
- $\langle 2 \rangle 6. \|T_m(u) T_n(u)\| < \epsilon$

Proof:

$$||T_m(u) - T_n(u)|| \le ||T_m - T_n|| ||u||$$
 (Proposition 18.2.11)

- $\langle 1 \rangle 3$ . Define  $T: U \to V$  by  $T(u) = \lim_{n \to \infty} T_n(u)$
- $\langle 1 \rangle 4. \ T \in \mathcal{B}(U, V)$ 
  - $\langle 2 \rangle 1$ . For all  $x, y \in U$  and  $\alpha, \beta \in K$  we have  $T(\alpha x + \beta y) = \alpha T(x) + \beta T(y)$ 
    - $\langle 3 \rangle 1$ . Let:  $x, y \in U$
    - $\langle 3 \rangle 2$ . Let:  $\alpha, \beta \in K$
    - $\langle 3 \rangle 3. \ T(\alpha x + \beta y) = \alpha T(x) + \beta T(y)$

Proof:

$$T(\alpha x + \beta y) = \lim_{n \to \infty} T_n(\alpha x + \beta y)$$
$$= \lim_{n \to \infty} (\alpha T_n(x) + \beta T_n(y))$$
$$= \alpha T(x) + \beta T(y)$$

- $\langle 2 \rangle 2$ . T is bounded.
  - $\langle 3 \rangle 1$ . PICK N such that  $\forall n \geq N . ||T_n T|| < 1$
  - $\langle 3 \rangle 2$ . Pick B > 0 such that  $\forall x \in U . ||T_N(x)|| \leq B||x||$
  - $\langle 3 \rangle 3$ . Let:  $x \in U$
  - $\langle 3 \rangle 4. \ \|T(x)\| \le (B+1)\|x\|$

Proof:

$$||T(x)|| \le ||T_N(x) - T(x)|| + ||T(x)||$$
 (Triangle inequality)  
 $\le ||T_N - T|| ||x|| + ||T|| ||x||$  (Proposition 18.2.11)  
 $< ||x|| + B||x||$  ( $\langle 3 \rangle 1, \langle 3 \rangle 2$ )  
 $= (B+1)||x||$ 

- $\langle 1 \rangle 5. \ T_n \to T \text{ as } n \to \infty$ 
  - $\langle 2 \rangle 1$ . Let:  $\epsilon > 0$
  - $\langle 2 \rangle 2$ . Pick N such that  $\forall m, n \geq N . ||T_m T_n|| < \epsilon/2$
  - $\langle 2 \rangle 3$ . Let:  $n \geq N$ Prove:  $||T_n - T|| < \epsilon$
  - $\langle 2 \rangle 4$ . Let:  $x \in U$  with ||x|| = 1
  - $\langle 2 \rangle 5$ .  $||T_n(x) T(x)|| < \epsilon/2$

PROOF: Let  $n \to \infty$  in  $\langle 2 \rangle 2$ .

**Corollary 18.3.17.1.** For any normed space V over K, the space  $\mathcal{B}(V,K)$  is a Banach space.

**Theorem 18.3.18.** Let U be a normed space and V a Banach space. Let W be a subspace of U. Let  $T:W\to V$  be a continuous linear transformation. Then T has a unique extension to a continuous linear transformation  $\operatorname{cl} W\to V$ .

#### Proof:

- $\langle 1 \rangle 1$ . Define  $S: \operatorname{cl} W \to V$  by:  $S(x) = \lim_{n \to \infty} T(x_n)$ , where  $(x_n)$  is any sequence in W that converges to x.
  - $\langle 2 \rangle 1$ . For all  $x \in \operatorname{cl} W$ , there exists a sequence  $(x_n)$  in W that converges to x. PROOF: Theorem 18.1.43.
  - $\langle 2 \rangle 2$ . If  $(x_n)$  is a Cauchy sequence in W then  $(T(x_n))$  is Cauchy in V.
    - $\langle 3 \rangle 1$ . Assume: w.l.o.g.  $T \neq 0$
    - $\langle 3 \rangle 2$ . Let:  $(x_n)$  be a Cauchy sequence in W.
    - $\langle 3 \rangle 3$ . Pick B > 0 such that  $\forall x \in W . ||T(x)|| \leq B||x||$
    - $\langle 3 \rangle 4$ . Let:  $\epsilon > 0$
    - $\langle 3 \rangle$ 5. PICK N such that  $\forall m, n \geq N . ||x_m x_n|| < \epsilon / ||T||$
    - $\langle 3 \rangle 6$ . Let:  $m, n \geq N$
    - $\langle 3 \rangle 7. \|T(x_m) T(x_n)\| < \epsilon$
  - (2)3. If  $(x_n)$  and  $(y_n)$  are sequences in W that converge to the same element in cl W then  $(T(x_n))$  and  $(T(y_n))$  have the same limit in V.
    - $\langle 3 \rangle 1$ . Assume: w.l.o.g.  $T \neq 0$
    - $\langle 3 \rangle 2$ . Assume:  $x_n \to l$  and  $y_n \to l$  as  $n \to \infty$
    - $\langle 3 \rangle 3$ . Let:  $T(x_n) \to a$  and  $T(y_n) \to b$  as  $n \to \infty$
    - $\langle 3 \rangle 4$ . Assume: for a contradiction  $a \neq b$
    - $\langle 3 \rangle 5$ . Let:  $\epsilon = ||a b||$
    - (3)6. PICK N such that  $\forall n \geq N. \|x_n l\| < \epsilon/3 \|T\|$  and  $\forall n \geq N. \|y_n l\| < \epsilon/3 \|T\|$
    - $\langle 3 \rangle 7. \ \forall n \geq N. ||T(x_n) T(y_n)|| < 2\epsilon/3$
    - $\langle 3 \rangle 8. \ \|a b\| \le 2\epsilon/3$
    - $\langle 3 \rangle 9$ . This contradicts  $\langle 3 \rangle 5$ .
- $\langle 1 \rangle 2$ . S extends T
  - $\langle 2 \rangle 1$ . Let:  $w \in W$
  - $\langle 2 \rangle 2$ .  $w \to w$  as  $n \to \infty$
  - $\langle 2 \rangle 3$ .  $T(w) \to T(w)$  as  $n \to \infty$
  - $\langle 2 \rangle 4$ . S(w) = T(w)
- $\langle 1 \rangle 3$ . S is bounded.
  - $\langle 2 \rangle 1$ . Let:  $x \in \operatorname{cl} W$

PROVE:  $||S(x)|| \le ||T|| ||x||$ 

- $\langle 2 \rangle 2$ . PICK a sequence  $(x_n)$  in W that converges to x.
- $\langle 2 \rangle 3$ .  $||T(x_n)|| \le ||T|| ||x_n||$  for all n.
- $\langle 2 \rangle 4. \ \| S(x) \| \le \| T \| \| x \|$

PROOF: Taking the limit as  $n \to \infty$ .

 $\langle 1 \rangle 4$ . S is linear.

- $\langle 2 \rangle 1$ . Let:  $x, y \in \operatorname{cl} W$  and  $\alpha, \beta \in K$
- $\langle 2 \rangle 2$ . PICK sequences  $(x_n)$  and  $(y_n)$  in W that converge to x and y.
- $\langle 2 \rangle 3$ .  $T(\alpha x_n + \beta y_n) = \alpha T(x_n) + \beta T(y_n)$  for all n.
- $\langle 2 \rangle 4$ .  $S(\alpha x + \beta y) = \alpha S(x) + \beta S(y)$

PROOF: Taking the limit as  $n \to \infty$ .

- $\langle 1 \rangle 5$ . S is unique.
  - $\langle 2 \rangle 1$ . Let: S' be a continuous linear extension of S defined on cl W.
  - $\langle 2 \rangle 2$ . Let:  $x \in W$ Prove: S(x) = S'(x)
  - $\langle 2 \rangle 3$ . PICK a sequence  $(x_n)$  in W that converges to x.
  - $\langle 2 \rangle 4$ .  $T(x_n) = S'(x_n) \to S'(x)$  as  $n \to \infty$
- $\langle 2 \rangle 5. \ S'(x) = S(x)$

Corollary 18.3.18.1. Let U be a normed space and V a Banach space. Let W be a dense subspace of U. Let  $T:W\to V$  be a continuous linear transformation. Then T has a unique extension to a continuous linear transformation  $U\to V$ .

**Definition 18.3.19** (Functional). Let V be a normed space over K. A functional on V is a bounded linear mapping  $V \to K$ . The dual space of V is the space  $\mathcal{B}(V,K)$  of all functionals.

**Theorem 18.3.20** (Banach-Steinhaus Theorem). Let  $\mathcal{T}$  be a family of bounded linear mappings from a Banach space X into a normed space Y. If, for every  $x \in X$ , there exists a constant  $M_x$  such that  $\forall T \in \mathcal{T}. ||T(x)|| \leq M_x$ , then there exists a constant M > 0 such that  $\forall T \in \mathcal{T}. ||T|| \leq M$ .

#### Proof:

- $\langle 1 \rangle 1$ . Assume: for a contradiction no such M exists.
- $\langle 1 \rangle 2$ . For  $n \in \mathbb{Z}_+$ , PICK  $T_n \in \mathcal{T}$  such that  $||T_n|| > n2^n$ .
- $\langle 1 \rangle 3$ . For  $n \in \mathbb{Z}_+$ , PICK  $x_n \in X$  such that  $||x_n|| = 1$  and  $||T_n(x_n)|| > n2^n$ .
- $\langle 1 \rangle 4$ . For  $n \in \mathbb{Z}_+$ ,

$$\left\| \frac{1}{n} T_n \left( \frac{x_n}{2^n} \right) \right\| > 1 .$$

- $\langle 1 \rangle 5$ . For  $i, j \in \mathbb{Z}_+$ , LET:  $y_{ij} = \frac{1}{i} T_i(\frac{x_j}{2^j})$ .
- $\langle 1 \rangle 6$ . For all  $j \in \mathbb{Z}_+$ ,  $y_{ij} \to 0$  as  $i \to \infty$ 
  - $\langle 2 \rangle 1$ . Let:  $j \in \mathbb{Z}_+$
  - $\langle 2 \rangle 2$ . Pick M such that  $\forall T \in \mathcal{T} . ||T(x_i/2^j)|| \leq M$
  - $\langle 2 \rangle 3$ . For all  $i, ||y_{ij}|| \leq M/i$
- $\langle 1 \rangle$ 7. For any increasing sequence of positive integers  $(p_i)$ , we have  $\sum_{j=1}^{\infty} y_{p_i p_j} \to 0$  as  $i \to \infty$ 
  - $\langle 2 \rangle 1$ . Let:  $(p_i)$  be an increasing sequence of positive integers.
  - $\langle 2 \rangle 2$ . Let:  $z = \sum_{j=1}^{\infty} x_{p_j}/2^{p_j}$

PROOF: This converges by Theorem 18.3.13.

- $\langle 2 \rangle 3$ . PICK C such that  $\forall T \in \mathcal{T} . ||T(z)|| \leq C$
- $\langle 2 \rangle 4$ . For all i,  $\|\sum_{j=1}^{\infty} y_{p_i p_j}\| \leq C/p_i$ .

PROOF: 
$$\left\|\sum_{j=1}^{\infty}y_{p_{i}p_{j}}\right\| = \left\|\sum_{j=1}^{\infty}\frac{1}{p_{i}}T_{p_{i}}\left(\frac{x_{p_{j}}}{2^{p_{j}}}\right)\right\| \qquad (\langle 1\rangle 5)$$

$$= \frac{1}{p_{i}}\left\|T_{p_{i}}\left(\sum_{j=1}^{\infty}\frac{x_{p_{j}}}{2^{p_{j}}}\right)\right\| \qquad (T_{p_{i}} \text{ continuous})$$

$$= \frac{1}{p_{i}}\left\|T_{p_{i}}(z)\right\| \qquad (\langle 2\rangle 2)$$

$$\leq \frac{C}{p_{i}} \qquad (\langle 2\rangle 3)$$

$$\langle 2\rangle 5. \sum_{j=1}^{\infty}y_{p_{i}p_{j}} \to 0 \text{ as } i \to \infty$$

$$\langle 1\rangle 8. \ y_{ii} \to 0 \text{ as } i \to \infty$$
PROOF: Diagonal Theorem,  $\langle 1\rangle 6$ ,  $\langle 1\rangle 7$ .
$$\langle 1\rangle 9. \ \text{Q.E.D.}$$

PROOF: Diagonal Theorem,  $\langle 1 \rangle 6$ ,  $\langle 1 \rangle 7$ .

PROOF:  $\langle 1 \rangle 4$  and  $\langle 1 \rangle 8$  form a contradiction. 

#### 18.4 Contraction Mappings

**Definition 18.4.1** (Contraction Mapping). Let E be a normed space over K. Let  $A \subseteq E$ . A function  $f: A \to E$  is a contraction (mapping) iff there exists a real  $\alpha$  such that  $0 < \alpha < 1$  and

$$\forall x, y \in A. ||f(x) - f(y)|| \le \alpha ||x - y||.$$

**Proposition 18.4.2.** Contraction mappings are uniformly continuous.

#### Proof:

- $\langle 1 \rangle 1$ . Let: E be a normed space over K.
- $\langle 1 \rangle 2$ . Let:  $A \subseteq E$
- $\langle 1 \rangle 3$ . Let:  $f: A \to E$  be a contraction mapping.
- $\langle 1 \rangle 4$ . PICK  $\alpha$  such that  $0 < \alpha < 1$  and  $\forall x, y \in A . || f(x) f(y) || \le \alpha || x y ||$ .
- $\langle 1 \rangle 5$ . Let:  $\epsilon > 0$
- $\langle 1 \rangle 6$ . Let:  $\delta = \epsilon / \alpha$
- $\langle 1 \rangle 7$ . For all  $x, y \in A$ , if  $||x y|| < \delta$  then  $||f(x) f(y)|| < \epsilon$ .

**Theorem 18.4.3** (Banach Fixed Point Theorem). Let E be a Banach space over K. Let F be a nonempty closed subset of E. Let  $f: F \to F$  be a contraction mapping. Then there exists a unique  $z \in F$  such that f(z) = z.

#### Proof:

 $\langle 1 \rangle 1$ . PICK  $\alpha$  such that  $0 < \alpha < 1$  and

$$\forall x, y \in F. ||f(x) - f(y)|| \le \alpha ||x - y||.$$

 $\langle 1 \rangle 2$ . Pick  $x_0 \in F$ 

$$\langle 1 \rangle 3$$
. For  $n \in \mathbb{Z}_+$ ,  
LET:  $x_n = f^n(x_0)$ .

- $\langle 1 \rangle 4$ .  $(x_n)$  is a Cauchy sequence.
  - $\langle 2 \rangle 1$ . For all  $n \in \mathbb{Z}_+$  we have  $||x_{n+1} x_n|| \le \alpha^n ||x_1 x_0||$ .
  - $\langle 2 \rangle 2$ . For all  $m, n \in \mathbb{Z}_+$  with m < n we have  $||x_n x_m|| < \alpha^m ||x_1 x_0||/(1-\alpha)$ .

Proof:

$$||x_{n} - x_{m}|| \leq ||x_{n} - x_{n-1}|| + ||x_{n-1} - x_{n-2}|| + \dots + ||x_{m+1} - x_{m}|| \quad \text{(Triangle inequality)}$$

$$\leq (\alpha^{n-1} + \alpha^{n-2} + \dots + \alpha^{m})||x_{1} - x_{0}||$$

$$< \frac{||x_{1} - x_{0}||}{1 - \alpha} \alpha^{m}$$

$$\langle 2 \rangle 3. \text{ Let: } \epsilon > 0$$

- $\langle 2 \rangle 4$ . PICK N such that  $\alpha^N ||x_1 x_0||/(1 \alpha) < \epsilon$
- $\langle 2 \rangle 5$ . For all  $m, n \geq N$ , we have  $||x_n x_m|| < \epsilon$
- $\langle 1 \rangle 5$ . Let:  $z = \lim_{n \to \infty} x_n$
- $\langle 1 \rangle 6. \ f(z) = z$

$$f(z) = f\left(\lim_{n \to \infty} x_n\right)$$

$$= \lim_{n \to \infty} f(x_n)$$
 (Proposition 18.4.2)
$$= \lim_{n \to \infty} x_{n+1}$$

$$= z$$

- $\langle 1 \rangle 7$ . For any  $w \in F$ , if f(w) = w then w = z.
  - $\langle 2 \rangle 1$ . Let:  $w \in F$
  - $\langle 2 \rangle 2$ . Assume: f(w) = w
  - $\langle 2 \rangle 3. \|z w\| \le \alpha \|z w\|$

PROOF: 
$$||z - w|| = ||f(z) - f(w)|| \le \alpha ||z - w||$$

- $\langle 2 \rangle 4. \ \|z w\| = 0$
- $\langle 2 \rangle 5. \ z = w$

### Chapter 19

# Inner Product Spaces

**Definition 19.0.1** (Inner Product Space). Let E be a complex vector space. An *inner product* on E is a function  $\langle \ , \ \rangle : E^2 \to \mathbb{C}$  such that, for all  $x,y,z \in E$  and  $\alpha,\beta \in \mathbb{C}$ , we have:

- 1.  $\langle x, y \rangle = \overline{\langle y, x \rangle}$
- 2.  $\langle \alpha x + \beta y, z \rangle = \alpha \langle x, z \rangle + \beta \langle y, z \rangle$
- 3.  $\langle x, x \rangle \geq 0$
- 4. If  $\langle x, x \rangle = 0$  then x = 0

An inner product space consists of a complex vector space V and an inner product on V.

**Proposition 19.0.2.** *Let* E *be an inner product space. For any*  $x \in E$ *, we have*  $\langle x, x \rangle$  *is real.* 

Proof: Since  $\langle x, x \rangle = \overline{\langle x, x \rangle}$ .  $\square$ 

Proposition 19.0.3.

$$\langle x,\alpha y+\beta z\rangle=\overline{\alpha}\langle x,y\rangle+\overline{\beta}\langle x,z\rangle$$

Proposition 19.0.4.

$$\langle 0, y \rangle = \langle x, 0 \rangle = 0$$

**Proposition 19.0.5.** The function  $\langle \vec{x}, \vec{y} \rangle = \sum_{i=1}^{n} x_i \overline{y_i}$  is an inner product on  $\mathbb{C}^n$ .

**Proposition 19.0.6.** The function  $\langle (x_n), (y_n) \rangle = \sum_{i=1}^{\infty} x_i \overline{y_i}$  is an inner product on  $l^2$ .

**Proposition 19.0.7.** The function  $\langle f, g \rangle = \int_a^b f(x) \overline{g(x)} dx$  is an inner product on  $\mathcal{C}([a, b])$ .

**Proposition 19.0.8.** Let p > 1 and  $\Omega \subseteq \mathbb{R}^N$ . Let  $L^p(\Omega)$  be the set of all functions  $f: \Omega \to \mathbb{C}$  such that  $|f|^p$  is Lebesgue integrable.

The function  $\langle f, g \rangle = \int_{\Omega} f(x) \overline{g(x)} dx$  is an inner product on  $L^2(\Omega)$ .

**Proposition 19.0.9.** Let  $E_1$  and  $E_2$  be inner product spaces. Then the function  $\langle (e_1, e_2), (e'_1, e'_2) \rangle = \langle e_1, e'_1 \rangle + \langle e_2, e'_2 \rangle$  is an inner product on  $E_1 \times E_2$ .

**Definition 19.0.10** (Norm). In an inner product space, define  $||x|| = \sqrt{\langle x, x \rangle}$ .

Proposition 19.0.11 (Schwarz's Inequality). In any inner product space,

$$|\langle x, y \rangle| \le ||x|| ||y||$$
.

Equality holds iff x and y are linearly dependent.

#### Proof:

- $\langle 1 \rangle 1$ . Assume: w.l.o.g.  $y \neq 0$
- $\langle 1 \rangle 2. \ |\langle x, y \rangle| \le ||x|| ||y||$ 
  - $\langle 2 \rangle 1$ . For all  $\alpha \in \mathbb{C}$  we have  $\langle x, x \rangle + \overline{\alpha} \langle x, y \rangle + \alpha \langle y, x \rangle + |\alpha|^2 \langle y, y \rangle$ PROOF: The right-hand side is  $\langle x + \alpha y, x + \alpha y \rangle$ .
  - $\langle 2 \rangle 2$ .  $\langle x, x \rangle \langle y, y \rangle |\langle x, y \rangle|^2 \ge 0$

Proof: Taking  $\alpha = -\langle x, x \rangle / \langle y, y \rangle$  in  $\langle 2 \rangle 1$ .

- $\langle 1 \rangle 3$ . If  $|\langle x, y \rangle| = ||x|| ||y||$  then x and y are linearly dependent.
  - $\langle 2 \rangle 1$ . Assume:  $|\langle x, y \rangle| = ||x|| ||y||$
  - $\langle 2 \rangle 2. \ \langle x, y \rangle \langle y, x \rangle = \langle x, x \rangle \langle y, y \rangle$
  - $\langle 2 \rangle 3. \ \langle y, y \rangle x \langle x, x \rangle y = 0$

Proof:

$$\langle \langle y, y \rangle x - \langle x, y \rangle y, \langle y, y \rangle x - \langle x, y \rangle y \rangle = \langle y, y \rangle^2 \langle x, x \rangle - \langle y, y \rangle \langle y, x \rangle \langle x, y \rangle - \langle x, y \rangle \langle y, y \rangle \langle y, x \rangle + \langle x, y \rangle \langle y, x \rangle = 0$$

- $\langle 1 \rangle 4$ . If x and y are linearly dependent then  $|\langle x, y \rangle| = ||x|| ||y||$ 
  - $\langle 2 \rangle 1$ . Assume: x and y are linearly dependent.
  - $\langle 2 \rangle 2$ . Let:  $y = \alpha x$
  - $\langle 2 \rangle 3. \ |\langle x, y \rangle| = ||x|| ||y||$

Proof:

$$\begin{aligned} |\langle x, y \rangle| &= |\langle x, \alpha x \rangle| \\ &= |\alpha| |\langle x, x \rangle| \\ &= |\alpha| ||x||^2 \\ &= ||x|| ||\alpha x|| \\ &= ||x|| ||y|| \end{aligned}$$

Corollary 19.0.11.1 (Triangle Inequality). In any inner product space,

$$||x + y|| \le ||x|| + ||y||$$

Proof:

$$||x + y||^2 = \langle x + y, x + y \rangle$$

$$= \langle x, x \rangle + \langle x, y \rangle + \langle y, x \rangle + \langle y, y \rangle$$

$$\leq \langle x, x \rangle + 2|\langle x, y \rangle| + \langle y, y \rangle$$

$$\leq ||x||^2 + 2||x|| ||y|| + ||y||^2 \qquad \text{(Schwarz's Inequality)}$$

$$= (||x|| + ||y||)^2 \qquad \Box$$

Corollary 19.0.11.2. The norm in an inner product space is a norm.

**Theorem 19.0.12** (Parallelogram Law). In any inner product space,

$$||x + y||^2 + ||x - y||^2 = 2(||x||^2 + ||y||^2)$$

Proof:

$$\begin{array}{ll} \langle 1 \rangle 1. \ \|x+y\|^2 = \|x\|^2 + \langle x,y \rangle + \langle y,x \rangle + \|y\|^2 \\ \langle 1 \rangle 2. \ \|x-y\|^2 = \|x\|^2 - \langle x,y \rangle - \langle y,x \rangle + \|y\|^2 \\ \langle 1 \rangle 3. \ \mathrm{Q.E.D.} \end{array}$$

PROOF: Add  $\langle 1 \rangle 1$  and  $\langle 1 \rangle 2$ .

**Proposition 19.0.13.** Let E be a normed space over  $\mathbb{C}$ . Then there exists an inner product on E that induces the norm of E iff E satisfies the Parallelogram Law.

Proof: If E satisfies the parallelogram law, define

$$\langle x, y \rangle = \frac{1}{4} (\|x + y\|^2 - \|x - y\|^2 + i\|x + iy\|^2 - i\|x - iy\|^2)$$
.

**Definition 19.0.14** (Orthogonal). Vectors x and y in an inner product space are orthogonal,  $x \perp y$ , iff  $\langle x, y \rangle = 0$ .

**Theorem 19.0.15** (Pythagorean Formula). If x and y are orthogonal then

$$||x+y||^2 = ||x||^2 + ||y||^2$$
.

**Definition 19.0.16** (Weak Convergence). Let E be an inner product space. Let  $(x_n)$  be a sequence in E and  $l \in E$ . Then  $(x_n)$  weakly converges to l,  $x_n \stackrel{w}{\to} l$ as  $n \to \infty$ , iff  $\forall y \in E.\langle x_n, y \rangle \to \langle l, y \rangle$  as  $n \to \infty$ .

**Proposition 19.0.17.** In any inner product space E, the inner product  $\langle , \rangle$ :  $E^2 \to \mathbb{C}$  is continuous.

$$\langle 1 \rangle 1$$
. Let:  $x_n \to x$  and  $y_n \to y$  in  $E$ .

$$\langle 1 \rangle 2. \ \langle x_n, y_n \rangle \to \langle x, y \rangle$$

Proof:

$$\begin{aligned} |\langle x_n, y_n \rangle - \langle x, y \rangle| &\leq |\langle x_n, y_n \rangle - \langle x_n, y \rangle| + |\langle x_n, y \rangle - \langle x, y \rangle| \\ &= |\langle x_n, y_n - y \rangle| + |\langle x_n - x, y \rangle| \\ &\leq \|x_n\| \|y_n - y\| + \|x_n - x\| \|y\| \end{aligned}$$
 (Schwarz's Inequality) 
$$\to 0$$

using the fact that  $(x_n)$  is bounded.

**Theorem 19.0.18.**  $x_n \to l$  if and only if  $x_n \stackrel{w}{\to} l$  and  $||x_n|| \to ||x||$ .

 $\langle 1 \rangle 1$ . If  $x_n \to l$  then  $x_n \stackrel{w}{\to} l$  and  $||x_n|| \to ||l||$ .

PROOF: Easy using the fact that the inner product is continuous.

- $\langle 1 \rangle 2$ . If  $x_n \stackrel{w}{\to} l$  and  $||x_n|| \to ||l||$  then  $x_n \to l$ .
  - $\langle 2 \rangle 1$ . Assume:  $x_n \stackrel{w}{\to} l$  and  $||x_n|| \to ||l||$  $\langle 2 \rangle 2$ .  $\langle x_n, l \rangle \to ||l||^2$

  - $\langle 2 \rangle 3. \|x_n l\| \to 0$

Proof:

$$||x_n - l||^2 = \langle x_n - l, x_n - l \rangle$$

$$= \langle x_n, x_n \rangle - \langle x_n, l \rangle - \langle l, x_n \rangle + \langle l, l \rangle$$

$$= ||x_n||^2 - \langle x_n, l \rangle - \overline{\langle x_n, l \rangle} + ||l||^2$$

$$\to ||l||^2 - 2||l||^2 + ||l||^2$$

$$= 0$$

**Theorem 19.0.19.** Let S be a subset of an inner product space E such that span S is dense in E. If  $(x_n)$  is a bounded sequence in E and, for all  $y \in S$ , we have  $\langle x_n, y \rangle \to \langle x, y \rangle$  then  $x_n \stackrel{w}{\to} x$ .

Proof:

- $\langle 1 \rangle 1$ . For all  $y \in \operatorname{span} S$ , we have  $\langle x_n, y \rangle \to \langle x, y \rangle$
- $\langle 1 \rangle 2$ . Let:  $z \in E$

Prove:  $\langle x_n, z \rangle \to \langle x, z \rangle$ 

 $\langle 1 \rangle 3$ . Let:  $\epsilon > 0$ 

PROVE: There exists  $n_0$  such that  $\forall n \geq n_0 . |\langle x_n, z \rangle - \langle x, z \rangle| < \epsilon$ 

- $\langle 1 \rangle 4$ . PICK M > 0 such that  $||x|| \leq M$  and  $\forall n \in \mathbb{Z}_+ . ||x_n|| \leq M$ .
- $\langle 1 \rangle 5$ . Pick  $y_0 \in \operatorname{span} S$  such that  $||z y_0|| < \epsilon/3M$
- $\langle 1 \rangle 6$ . Pick  $n_0 \in \mathbb{Z}_+$  such that, for all  $n \geq n_0$ , we have  $|\langle x_n, y_0 \rangle \langle x, y_0 \rangle| < \epsilon/3$
- $\langle 1 \rangle 7$ . Let:  $n \geq n_0$
- $\langle 1 \rangle 8. \ |\langle x_n, z \rangle \langle x, z \rangle| < \epsilon$

Proof:

$$\begin{split} |\langle x_n, z \rangle - \langle x, z \rangle| &\leq |\langle x_n, z \rangle - \langle x_n, y_0 \rangle| + |\langle x_n, y_0 \rangle - \langle x, y_0 \rangle| + |\langle x, y_0 \rangle - \langle x, z \rangle| \\ &< \|x_n\| \|z - y_0\| + \epsilon/3 + \|x\| \|y_0 - z\| \\ &< M(\epsilon/3M) + \epsilon/3 + M(\epsilon/3M) \\ &= \epsilon \end{split}$$

#### 19.1 Orthonormal Bases

**Definition 19.1.1** (Orthogonal). Let V be an inner product space and  $S \subseteq V$ . Then S is *orthogonal* iff any two distinct elements of S are orthogonal.

**Definition 19.1.2** (Orthonormal). Let V be an inner product space and  $S \subseteq V$ . Then S is orthonormal iff it is orthogonal and  $\forall x \in S. ||x|| = 1$ .

Proposition 19.1.3. Orthonormal sets are linearly independent.

### Proof:

 $\langle 1 \rangle 1$ . Let: S be orthonormal

 $\langle 1 \rangle 2$ . Assume:  $\alpha_1 e_1 + \cdots + \alpha_n e_n = 0$  where  $e_1, \dots, e_n \in S$   $\langle 1 \rangle 3$ .  $|\alpha_1|^2 + \cdots + |\alpha_n|^2 = 0$ 

$$\langle 1 \rangle 3. \ |\alpha_1|^2 + \dots + |\alpha_n|^2 = 0$$

Proof:

$$0 = \sum_{m=1}^{n} \langle 0, \alpha_m e_m \rangle$$

$$= \sum_{m=1}^{n} \langle \sum_{k=1}^{n} \alpha_k e_k, \alpha_m e_m \rangle$$

$$= \sum_{m=1}^{n} \sum_{k=1}^{n} \alpha_k \overline{\alpha_m} \langle e_k, e_m \rangle$$

$$= \sum_{k=1}^{n} |\alpha_k|^2$$

$$\langle 1 \rangle 4. \ \alpha_1 = \dots = \alpha_n = 0$$

**Proposition 19.1.4.** In  $l^2$ , let  $e_n$  be the sequence whose nth element is 1 and whose other elements are 0. Then  $\{e_n \mid n \in \mathbb{Z}_+\}$  is orthonormal.

**Proposition 19.1.5.** In  $L^2([-\pi,\pi])$ , let  $\phi_n(x) = e^{inx}/\sqrt{2\pi}$  for  $n \in \mathbb{Z}$ . Then  $\{\phi_n \mid n \in \mathbb{Z}\}\ is\ orthonormal.$ 

**Definition 19.1.6** (Legendre Polynomials). The Legendre polynomials  $P_n \in$  $\mathbb{Q}[x]$  for  $n \in \mathbb{N}$  are defined by

$$P_0 = 1$$

$$P_n = \frac{1}{2^n n!} \frac{d^n}{dx^n} (x^2 - 1)^n$$

**Proposition 19.1.7.** Let  $P_n$  be the nth Legendre polynomial. Then  $\{P_n \mid n \in P_n \mid n \in P_n$  $\mathbb{N}$  is orthogonal in  $L^2([-1,1])$ .

**Definition 19.1.8** (Hermite Polynomial). The Hermite polynomials  $H_n \in \mathbb{R}[x]$ for  $n \in \mathbb{N}$  are defined by

$$H_n(x) = (-1)^n e^{x^2} \frac{d^n}{dx^n} e^{-x^2}$$
.

**Proposition 19.1.9.** Let  $H_n$  be the nth Hermite polynomial. Then  $\{e^{-x^2/2}H_n(x)\mid$  $n \in \mathbb{N}$  is orthogonal in  $L^2(\mathbb{R})$ .

**Theorem 19.1.10.** Let V be an inner product space. If  $x_1, \ldots, x_n \in V$  are orthogonal then

$$\left\| \sum_{i=1}^{n} x_i \right\|^2 = \sum_{i=1}^{n} \|x_i\|^2.$$

**Theorem 19.1.11** (Bessel's Equality). Let V be an inner product space. Let  $x_1, \ldots, x_n \in V$  be orthonormal. Let  $x \in V$ . Then

$$\left\| x - \sum_{k=1}^{n} \langle x, x_k \rangle x_k \right\|^2 = \|x\|^2 - \sum_{k=1}^{n} |\langle x, x_k \rangle|^2.$$

Proof.

$$\left\| x - \sum_{k=1}^{n} \langle x, x_k \rangle x_k \right\|^2 = \left\langle x - \sum_{k=1}^{n} \langle x, x_k \rangle x_k, x - \sum_{k=1}^{n} \langle x, x_k \rangle x_k \right\rangle$$

$$= \langle x, x \rangle - \left\langle x, \sum_{k=1}^{n} \langle x, x_k \rangle x_k \right\rangle - \left\langle \sum_{k=1}^{n} \langle x, x_k \rangle x_k, x \right\rangle$$

$$+ \left\langle \sum_{k=1}^{n} \langle x, x_k \rangle x_k, \sum_{k=1}^{n} \langle x, x_k \rangle x_k \right\rangle$$

$$= \langle x, x \rangle - 2 \sum_{k=1}^{n} \langle x, x_k \rangle \langle x_k, x \rangle + \sum_{i=1}^{n} \sum_{j=1}^{n} \langle x, x_i \rangle \langle x_j, x \rangle \langle x_i, x_j \rangle$$

$$= \|x\|^2 - 2 \sum_{k=1}^{n} |\langle x, x_k \rangle|^2 + \sum_{i=1}^{n} \langle x, x_i \rangle \langle x_i, x \rangle$$

$$= \|x\|^2 - \sum_{k=1}^{n} |\langle x, x_k \rangle|^2$$

**Corollary 19.1.11.1** (Bessel's Inequality). Let V be an inner product space. Let  $x_1, \ldots, x_n \in V$  be orthonormal. Let  $x \in E$ . Then

$$\sum_{k=1}^{n} |\langle x, x_k \rangle|^2 \le ||x||^2.$$

Corollary 19.1.11.2. Orthonormal sequences are weakly convergent to 0.

PROOF: Let  $(x_n)$  be an orthonormal sequence. Taking the limit in Bessel's inequality we have  $\sum_{k=1}^{\infty} |\langle x, x_k \rangle|^2 \leq ||x||^2 < \infty$  and so  $\langle x, x_k \rangle \to 0$  as  $k \to \infty$ .

**Corollary 19.1.11.3** (Generalized Fourier Series). Let V be an inner product space. Let  $(e_n)$  be an orthonormal sequence in V. For any  $x \in V$ , the generalized Fourier series of x is

$$\sum_{n=1}^{\infty} \langle x, e_n \rangle e_n ,$$

and  $\langle x, e_n \rangle$  is called the nth generalized Fourier coefficient of x with respect to  $(e_n)$ . We have  $(\langle x, e_n \rangle e_n)_n \in l^2$ .

**Definition 19.1.12** (Complete Orthonormal Sequence). Let E be an inner product space. Let  $(x_n)$  be an orthonormal sequence in E. Then  $(x_n)$  is *complete* iff, for all  $x \in E$ , we have

$$\sum_{n=1}^{\infty} \langle x, x_n \rangle x_n = x .$$

## Chapter 20

# Hilbert Spaces

**Definition 20.0.1** (Hilbert Space). A *Hilbert space* is a complete inner product space.

**Proposition 20.0.2.** For  $n \in \mathbb{N}$ ,  $\mathbb{C}^n$  is a Hilbert space.

**Proposition 20.0.3.**  $l^2$  is a Hilbert space.

**Proposition 20.0.4.**  $L^2(\mathbb{R})$  is a Hilbert space.

**Proposition 20.0.5.**  $L^2([a,b])$  is a Hilbert space.

**Proposition 20.0.6.** Let  $\rho$  be a measurable function on [a,b] such that  $\rho(x) > 0$  almost everywhere. Let  $L^{2\rho}([a,b])$  be the set of all measurable functions  $f:[a,b] \to \mathbb{C}$  such that

$$\int_{a}^{b} |f(x)|^{2} \rho(x) dx < \infty .$$

Define an inner product on  $L^{2\rho}([a,b])$  by

$$\langle f,g\rangle = \int_a^b f(x)\overline{g(x)}\rho(x)dx \ .$$

Then  $L^{2\rho}([a,b])$  is a Hilbert space.

**Proposition 20.0.7.** Let m and N be positive integers. Let  $\Omega$  be an open set in  $\mathbb{R}^N$ . Let  $\tilde{H}^m(\Omega)$  be the set of all  $f \in \mathcal{C}^m(\Omega)$  such that, for every  $\alpha = (\alpha_1, \ldots, \alpha_N) \in \mathbb{Z}_+^N$  with  $|\alpha| := \alpha_1 + \cdots + \alpha_N \leq m$ , we have

$$D^{\alpha}f := \frac{\partial^{|\alpha|}f}{\partial x_1^{\alpha_1}\partial x_2^{\alpha_2}\cdots\partial x_N^{\alpha_N}} \in L^2(\Omega) .$$

Define an inner product on  $\tilde{H}^m(\Omega)$  by

$$\langle f, g \rangle := \int_{\Omega} \sum_{\alpha} D^{\alpha} f \overline{D^{\alpha} g} .$$

Let  $H^m(\Omega)$  be the completion of  $\tilde{H}^m(\Omega)$ . Then  $H^m(\Omega)$  is a Hilbert space.

**Theorem 20.0.8.** Weakly convergent sequences in a Hilbert space are bounded.

 $\langle 1 \rangle 1$ . Let: H be a Hilbert space.

 $\langle 1 \rangle 2$ . Let:  $(x_n)$  be a weakly convergent sequence in H.

 $\langle 1 \rangle 3$ . For  $n \in \mathbb{Z}_+$ , Let:  $f_n: H \to \mathbb{C}, f_n(x) = \langle x, x_n \rangle$ 

 $\langle 1 \rangle 4$ . For  $n \in \mathbb{Z}_+$ ,  $f_n$  is a bounded linear functional.

 $\langle 1 \rangle 5$ . For every  $x \in H$ , the sequence  $(f_n(x))$  is bounded.

PROOF: Since it converges.

 $\langle 1 \rangle 6$ . Pick M > 0 such that, for all  $n \in \mathbb{Z}_+$ , we have  $||f_n|| \leq M$ . PROOF: Banach-Steinhaus Theorem,  $\langle 1 \rangle 4$ ,  $\langle 1 \rangle 5$ .

 $\langle 1 \rangle 7. \ \forall n \in \mathbb{Z}_+. ||f_n|| = ||x_n||$ 

 $\langle 2 \rangle 1$ . Let:  $n \in \mathbb{Z}_+$ 

 $\langle 2 \rangle 2$ .  $||f_n|| \leq ||x_n||$ 

PROOF: Since for all  $x \in H$  we have  $|f_n(x)| = |\langle x, x_n \rangle| \leq ||x|| ||x_n||$  by Schwarz's Inequality.

 $\langle 2 \rangle 3$ .  $||x_n|| \leq ||f_n||$ 

PROOF: Since  $||x_n||^2 = |\langle x_n, x_n \rangle| = |f_n(x_n)| \le ||f_n|| ||x_n||$ .

 $\langle 1 \rangle 8. \ \forall n \in \mathbb{Z}_+. ||x_n|| \leq M$ 

Proof:  $\langle 1 \rangle 6$ ,  $\langle 1 \rangle 7$ 

**Theorem 20.0.9.** Let H be a Hilbert space. Let  $(x_n)$  be an orthonormal sequence in H and let  $(\alpha_n)$  be a sequence of complex numbers. Then the series  $\sum_{n=1}^{\infty} \alpha_n x_n$  converges in H if and only if  $\sum_{n=1}^{\infty} |\alpha_n|$  converges in  $\mathbb{R}$ , in which

$$\left\| \sum_{n=1}^{\infty} \alpha_n x_n \right\|^2 = \sum_{n=1}^{\infty} |\alpha_n|^2.$$

PROOF:

 $\langle 1 \rangle 1$ . For m > k > 0 we have

$$\left\| \sum_{n=k}^{m} \alpha_n x_n \right\|^2 = \sum_{n=k}^{m} |\alpha_n|^2.$$

PROOF: Theorem 19.1.10.

 $\langle 1 \rangle 2$ . If  $\sum_{n=1}^{\infty} |\alpha_n|^2 < \infty$  then  $\sum_{n=1}^{\infty} \alpha_n x_n$  converges.  $\langle 2 \rangle 1$ . Assume:  $\sum_{n=1}^{\infty} |\alpha_n|^2 < \infty$ 

 $\langle 2 \rangle 2$ .  $(\sum_{n=1}^{m} \alpha_n x_n)_m$  is Cauchy. PROOF: From  $\langle 1 \rangle 1$ .

 $\langle 2 \rangle 3. \sum_{n=1}^{\infty} \alpha_n x_n$  converges.  $\langle 1 \rangle 3. \text{ If } \sum_{n=1}^{\infty} \alpha_n x_n$  converges then  $\sum_{n=1}^{\infty} |\alpha_n|^2 < \infty$ .

PROOF: From  $\langle 1 \rangle 1$ .  $\langle 1 \rangle 4$ . If  $\sum_{n=1}^{\infty} |\alpha_n|^2 < \infty$  then

$$\left\| \sum_{n=1}^{\infty} \alpha_n x_n \right\|^2 = \sum_{n=1}^{\infty} |\alpha_n|^2.$$

Proof: From  $\langle 1 \rangle 1$ .

**Proposition 20.0.10.** Every complete orthonormal sequence in a Hilbert space is a basis.

### Proof:

- $\langle 1 \rangle 1$ . Let: E be an inner product space.
- $\langle 1 \rangle 2$ . Let:  $(e_n)$  be a complete orthonormal sequence in E.
- $\langle 1 \rangle 3$ . For all  $x \in E$ , there exists a sequence  $(\alpha_n)$  in  $\mathbb{C}$  such that  $x = \sum_n \alpha_n e_n$ . PROOF: Immediate from  $\langle 1 \rangle 2$ .
- $\langle 1 \rangle 4$ . If  $\sum_{n} \alpha_{n} e_{n} = \sum_{n} \beta_{n} e_{n}$  then  $\alpha_{n} = \beta_{n}$  for all n.  $\langle 2 \rangle 1$ . Let:  $x = \sum_{n} \alpha_{n} e_{n} = \sum_{n} \beta_{n} e_{n}$   $\langle 2 \rangle 2$ .  $\sum_{n} |\alpha_{n} \beta_{n}|^{2} = 0$

Proof:

$$0 = \|x - x\|^{2}$$

$$= \left\| \sum_{n=1}^{\infty} \alpha_{n} e_{n} - \sum_{n=1}^{\infty} \beta_{n} e_{n} \right\|^{2}$$

$$= \left\| \sum_{n=1}^{\infty} (\alpha_{n} - \beta_{n}) e_{n} \right\|^{2}$$

$$= \sum_{n=1}^{\infty} |\alpha_{n} - \beta_{n}|^{2}$$
(Theorem 20.0.9)

 $\langle 2 \rangle 3$ .  $\alpha_n = \beta_n$  for all n.

**Theorem 20.0.11.** An orthonormal sequence  $(x_n)$  in a Hilbert space H is complete if and only if, for all  $x \in H$ , if  $\forall n. \langle x, x_n \rangle = 0$  then x = 0.

### Proof:

- $\langle 1 \rangle 1$ . If  $(x_n)$  is complete then, for all  $x \in H$ , if  $\forall n. \langle x, x_n \rangle = 0$  then x = 0.
  - $\langle 2 \rangle 1$ . Assume:  $(x_n)$  is complete.
  - $\langle 2 \rangle 2$ . Let:  $x \in H$
- $\langle 2 \rangle 3$ . Assume:  $\forall n. \langle x, x_n \rangle = 0$  $\langle 2 \rangle 4$ .  $x = \sum_{n=1}^{\infty} \langle x, x_n \rangle x_n = 0$  $\langle 1 \rangle 2$ . If, for all  $x \in H$ , if  $\forall n. \langle x, x_n \rangle = 0$  then x = 0, then  $(x_n)$  is complete.
  - $\langle 2 \rangle 1$ . Assume: For all  $x \in H$ , if  $\forall n. \langle x, x_n \rangle = 0$ , then x = 0.  $\langle 2 \rangle 2$ . Let:  $y = x \sum_{n=1}^{\infty} \langle x, x_n \rangle x_n$   $\langle 2 \rangle 3$ . For all  $n, \langle y, x_n \rangle = 0$

  - - $\langle 3 \rangle 1$ . Let:  $n \in \mathbb{Z}_+$
    - $\langle 3 \rangle 2. \ \langle y, x_n \rangle = 0$

Proof:

$$\langle y, x_n \rangle = \left\langle x - \sum_{m=1}^{\infty} \langle x, x_m \rangle x_m, x_n \right\rangle$$
$$= \langle x, x_n \rangle - \sum_{m=1}^{\infty} \langle x, x_m \rangle \langle x_m, x_n \rangle$$
$$= \langle x, x_n \rangle - \langle x, x_n \rangle$$
$$= 0$$

$$\langle 2 \rangle 4. \ y = 0$$
  
 $\langle 2 \rangle 5. \ x = \sum_{n=1}^{\infty} \langle x, x_n \rangle x_n$ 

**Theorem 20.0.12** (Parseval's Formula). Let H be a Hilbert space. Let  $(x_n)$ be an orthonormal sequence in H. Then  $(x_n)$  is complete if and only if, for all  $x \in H$ ,

$$||x||^2 = \sum_{n=1}^{\infty} |\langle x, x_n \rangle|^2.$$

- $\langle 1 \rangle 1$ . If  $(x_n)$  is complete then for all  $x \in H$  we have  $||x||^2 = \sum_{n=1}^{\infty} |\langle x, x_n \rangle|^2$ .
  - $\langle 2 \rangle 1$ . Assume:  $(x_n)$  is complete.

  - $\langle 2 \rangle 2$ . Let:  $x \in H$   $\langle 2 \rangle 3$ .  $||x||^2 = \sum_{n=1}^{\infty} |\langle x, x_n \rangle|^2$ PROOF:

$$||x||^2 = \left\| \sum_{n=1}^{\infty} \langle x, x_n \rangle x_n \right\|^2$$

$$= \sum_{n=1}^{\infty} |\langle x, x_n \rangle|^2$$
(Theorem 20.0.9)

- $\langle 1 \rangle 2$ . If, for all  $x \in H$ , we have  $||x||^2 = \sum_{n=1}^{\infty} |\langle x, x_n \rangle|^2$ , then  $(x_n)$  is complete.  $\langle 2 \rangle 1$ . Assume: For all  $x \in H$ , we have  $||x||^2 = \sum_{n=1}^{\infty} |\langle x, x_n \rangle|^2$
- $\langle 2 \rangle 2$ . Let:  $x \in H$   $\langle 2 \rangle 3$ .  $x = \sum_{n=1}^{\infty} \langle x, x_n \rangle x_n$

**Proposition 20.0.13.** For  $n \in \mathbb{Z}$ , let  $\pi_n(x) = e^{inx}/\sqrt{2\pi}$ . Then  $\{\pi_n \mid n \in \mathbb{Z}\}$ is a complete orthonormal set in  $L^2([-\pi, \pi])$ .

TODO

**Proposition 20.0.14.**  $B = \{1/\sqrt{2\pi}\} \cup \{\cos nx/\sqrt{\pi} \mid n \in \mathbb{Z}_+\} \cup \{\sin nx/\sqrt$  $n \in \mathbb{Z}_+$  is a complete orthonormal set in  $L^2([-\pi, \pi])$ .

Proof:

 $\langle 1 \rangle 1$ . For all  $f \in B$  we have ||f|| = 1 $\langle 2 \rangle 1. \ \|1/\sqrt{2\pi}\| = 1$ 

Proof:

$$||1/\sqrt{2\pi}|| = \int_{-\pi}^{\pi} dx/2\pi$$

 $\langle 2 \rangle 2$ . For all  $n \in \mathbb{Z}_+$  we have  $\|\cos nx/\sqrt{\pi}\| = 1$  Proof:

$$\|\cos nx/\sqrt{\pi}\| = 1/\pi \int_{-\pi}^{\pi} \cos^2 nx \, dx$$

$$= 1/2\pi \int_{-\pi}^{\pi} (\cos 2nx + 1) dx$$

$$= 1/2\pi \left[ 1/2n \sin 2nx + x \right]_{-\pi}^{\pi}$$

$$= (1/2\pi)(2\pi)$$

$$= 1$$

 $\langle 2 \rangle 3$ . For all  $n \in \mathbb{Z}_+$  we have  $\|\sin nx/\sqrt{\pi}\| = 1$  PROOF:

$$\|\sin nx/\sqrt{\pi}\| = 1/\pi \int_{-\pi}^{\pi} \sin^2 nx \, dx$$

$$= -1/2\pi \int_{-\pi}^{\pi} (\cos 2nx - 1) dx$$

$$= -1/2\pi \left[ 1/2n \sin 2nx - x \right]_{-\pi}^{\pi}$$

$$= (-1/2\pi)(-2\pi)$$

$$= 1$$

 $\langle 1 \rangle 2.$  For all  $f,g \in B$  with  $f \neq g$  we have  $\langle f,g \rangle = 0$ 

 $\langle 2 \rangle 1. \ \langle 1, \cos nx \rangle = 0$ PROOF:

$$\int_{-\pi}^{\pi} \cos nx \, dx = \left[1/n \sin nx\right]_{-\pi}^{\pi}$$

 $\langle 2 \rangle 2$ .  $\langle 1, \sin nx \rangle = 0$ 

Proof:

$$\int_{-\pi}^{\pi} \sin nx \, dx = \left[ -1/n \cos nx \right]_{-\pi}^{\pi}$$
$$= -1/n \cos n\pi + 1/n \cos n\pi$$
$$= 0$$

 $\langle 2 \rangle 3$ . If  $m \neq n$  then  $\langle \cos mx, \cos nx \rangle = 0$ 

PROOF:  

$$\int_{-\pi}^{\pi} \cos mx \cos nx \, dx = 1/2 \int_{-\pi}^{\pi} (\cos(n+m)x - \cos(n-m)x) dx$$

$$= 1/2 \left[ \frac{1}{n+m} \sin(n+m)x - \frac{1}{n-m} \sin(n-m)x \right]_{-\pi}^{\pi}$$

$$= 0$$

 $\langle 2 \rangle 4. \ \langle \cos mx, \sin nx \rangle = 0$ 

PROOF: 
$$\int_{-\pi}^{\pi} \cos mx \sin nx \, dx = 1/2 \int_{-\pi}^{\pi} (\sin(n+m)x - \sin(n-m)x) dx$$
$$= 1/2 \left[ -\frac{1}{n+m} \cos(n+m)x + \frac{1}{n-m} \cos(n-m)x \right]_{-\pi}^{\pi}$$
$$= 0$$
 (cos is odd)

 $\langle 2 \rangle 5$ . If  $m \neq n$  then  $\langle \sin mx, \sin nx \rangle = 0$ 

PROOF:  

$$\int_{-\pi}^{\pi} \sin mx \sin nx \, dx = 1/2 \int_{-\pi}^{\pi} (\cos(n-m)x - \cos(n+m)x) dx$$

$$= 1/2 \left[ \frac{1}{n-m} \sin(n-m)x - \frac{1}{n+m} \sin(n+m)x \right]_{-\pi}^{\pi}$$

- $\langle 1 \rangle 3.$  For all  $f \in L^2([-\pi,\pi]),$  if  $\forall g \in B. \langle f,g \rangle = 0$  then f=0  $\langle 2 \rangle 1.$  Let:  $f \in L^2([-\pi,\pi])$ 

  - $\langle 2 \rangle 2$ . Assume:  $\forall g \in B. \langle f, g \rangle = 0$

 $\langle 2 \rangle 3$ . For all  $n \in \mathbb{Z}$ ,  $\langle f, e^{inx} \rangle = 0$ PROOF: Since  $e^{inx} = \cos nx + i \sin nx$ .

 $\langle 2 \rangle 4$ . f = 0

Proof: From Proposition 20.0.13.

**Proposition 20.0.15.**  $\{\frac{1}{\sqrt{\pi}}\}\cup\{\sqrt{\frac{2}{\pi}}\cos nx\mid n\in\mathbb{Z}_+\}\ is\ a\ complete\ orthonormal$ set in  $L^{2}([0,\pi])$ .

**Proposition 20.0.16.**  $\{\sqrt{\frac{2}{\pi}}\sin nx \mid n \in \mathbb{Z}_+\}\ is\ a\ complete\ orthonormal\ set\ in$  $L^2([0,\pi]).$ 

**Definition 20.0.17** (Signum). The *signum* function sgn :  $\mathbb{R} \to \mathbb{R}$  is defined by

$$\operatorname{sgn} x = \begin{cases} 1 & \text{if } x > 0 \\ 0 & \text{if } x = 0 \\ -1 & \text{if } x < 0 \end{cases}$$

**Definition 20.0.18** (Rademacher Functions). The Rademarcher functions R:  $\mathbb{N} \times [0,1] \to \{-1,0,1\}$  are defined by

$$R(m,x) = \operatorname{sgn}(\sin(2^m \pi x)) .$$

**Proposition 20.0.19.** The Rademacher functios  $\{R(m, -) \mid m \in \mathbb{N}\}$  are orthonormal in  $L^2([0,1])$ .

Proof:

 $\langle 1 \rangle 1. \ \forall m \in \mathbb{N}. ||R(m, -)|| = 1$ 

PROOF:  $\int_0^1 \operatorname{sgn}(\sin(2^m \pi x))^2 dx = 1$  since the integrand is 1 except for finitely many points in [0,1].

- $\langle 1 \rangle 2$ . Given natural numbers  $m \neq n$ , we have  $\langle R(m,-), R(n,-) \rangle = 0$ 
  - $\langle 2 \rangle 1$ . Given reals a, b and a natural number m, we have  $\int_a^b R(m,x)dx = 0$ whenever  $2^m(b-a)$  is an even integer.

PROOF: If m > 0, or if m = 0 and b - a is an even integer, then the regions where R(m, x) = 1 are isometric with the regions where R(m, x) = -1.

- $\langle 2 \rangle 2$ . Let: m and n be natural numbers with n < m.
- $\langle 2 \rangle 3. \langle R(m,-), R(n,-) \rangle = 0$

Proof:

$$\int_{0}^{1} R(m,x)R(n,x)dx = \sum_{k=1}^{2^{n}} \int_{\frac{k-1}{2^{n}}}^{\frac{k}{2^{n}}} R(m,x)R(n,x)dx$$

$$= \sum_{k=1}^{2^{n}} (-i)^{k+1} \int_{\frac{k-1}{2^{n}}}^{\frac{k}{2^{n}}} R(m,x)dx$$

$$= 0 \qquad (\langle 2 \rangle 1, 2^{m} \left(\frac{k}{2^{n}} - \frac{k-1}{2^{n}}\right) = 2^{m-n} \text{ is an even integer})$$

**Proposition 20.0.20.** The set of Rademacher functions is not complete.

Proof:

- (1)1. Define  $f:[0,1]\to\mathbb{C}$  by f(x)=0 if  $0\leq x<1/4$ , f(x)=1 if  $1/4\leq x\leq 1$ 3/4, f(x) = 0 if  $3/4 < x \le 1$ .
- $\langle 1 \rangle 2. \ f \in L^2([0,1])$
- $\langle 1 \rangle 3. \ \langle R(0, -), f \rangle = 1/2$
- $\langle 1 \rangle 4$ .  $\langle R(m, -), f \rangle = 0$  for  $m \ge 1$
- $\langle 1 \rangle 5. \ f \neq 1/2R(0, -)$

**Definition 20.0.21** (Walsh Functions). Define the Walsh functions  $W: \mathbb{N} \times \mathbb{N}$  $[0,1] \to \{-1,0,1\}$  as follows. Given  $m \in \mathbb{N}$ , let  $m = \sum_{k=1}^n 2^{k-1} a_k$  where each  $a_k$  is either 0 or 1. Then

$$W(m,x) = \prod_{k=1}^{n} R(k,x)^{a_k}$$
.

**Proposition 20.0.22.** The set of Walsh functions  $\{W(m,-) \mid m \in \mathbb{N}\}$  is a compete orthonormal set.

TODO