# Mathematics

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# Chapter 1

# Sets and Classes

## 1.1 Classes

Our language is the language of first-order logic with equality over one primitive binary predicate  $\in$ . We call all the objects we reason about *sets*. When  $a \in b$ , we say a is a *member* or *element* of b, or b contains a. We write  $b \ni a$  for  $a \in b$ , and  $a \notin b$  for  $\neg(a \in b)$ . We write  $\forall x \in a.\phi$  as an abbreviation for  $\forall x(x \in a \to \phi)$ , and  $\exists x \in a.\phi$  as an abbreviation for  $\exists x(x \in a \land \phi)$ .

We shall speak informally of *classes* as an abbreviation for talking about predicates. A *class* is determined by a unary predicate  $\phi[x]$  (possibly with parameters). We write  $\{x \mid \phi[x]\}$  or  $\{x : \phi[x]\}$  for the class determined by  $\phi[x]$ . We write 'a is an element of  $\{x \mid \phi[x]\}$ ' or ' $a \in \{x \mid \phi[x]\}$ ' for  $\phi[a]$ .

We write  $\{t[x_1, ..., x_n] \mid P[x_1, ..., x_n]\}$  for

$$\{y \mid \exists x_1, \dots, x_n (y = t[x_1, \dots, x_n] \land P[x_1, \dots, x_n])\}$$
.

We say two classes **A** and **B** are *equal*, and write  $\mathbf{A} = \mathbf{B}$ , iff  $\forall x (x \in \mathbf{A} \leftrightarrow x \in \mathbf{B})$ .

**Proposition Schema 1.1.1.** For any class **A**, the following is a theorem.

$$\mathbf{A} = \mathbf{A}$$

PROOF: We have  $\forall x (x \in \mathbf{A} \Leftrightarrow x \in \mathbf{A})$ .  $\square$ 

**Proposition Schema 1.1.2.** For any classes **A** and **B**, the following is a theorem.

If 
$$\mathbf{A} = \mathbf{B}$$
 then  $\mathbf{B} = \mathbf{A}$ .

PROOF: If  $\forall x (x \in \mathbf{A} \Leftrightarrow x \in \mathbf{B})$  then  $\forall x (x \in \mathbf{B} \Leftrightarrow x \in \mathbf{A})$ .

**Proposition Schema 1.1.3.** For any classes A, B and C, the following is a theorem.

If 
$$A = B$$
 and  $B = C$  then  $A = C$ .

PROOF: If  $\forall x (x \in \mathbf{A} \Leftrightarrow x \in \mathbf{B})$  and  $\forall x (x \in \mathbf{B} \Leftrightarrow x \in \mathbf{C})$  then  $\forall x (x \in \mathbf{A} \Leftrightarrow x \in \mathbf{C})$ .  $\Box$ 

### 1.1.1 Subclasses

**Definition 1.1.4** (Subclass). We say a class **A** is a *subclass* of **B**, or **B** is a *superclass* of **A**, or **B** *includes* **A**, and write  $\mathbf{A} \subseteq \mathbf{B}$  or  $\mathbf{B} \supseteq \mathbf{A}$ , iff every element of **A** is an element of **B**. Otherwise we write  $\mathbf{A} \not\subseteq \mathbf{B}$  or  $\mathbf{B} \not\supseteq \mathbf{A}$ .

We say **A** is a *proper* subclass of **B**, **B** is a *proper* superclass of **A**, or **B** properly includes **A**, and write  $\mathbf{A} \subsetneq \mathbf{B}$  or  $\mathbf{B} \supsetneq \mathbf{A}$ , iff  $\mathbf{A} \subseteq \mathbf{B}$  and  $\mathbf{A} \ne \mathbf{B}$ .

**Proposition Schema 1.1.5.** For any class A, the following is a theorem.

$$\mathbf{A} \subseteq \mathbf{A}$$

PROOF: Every element of **A** is an element of **A**.  $\square$ 

**Proposition Schema 1.1.6.** For any classes **A** and **B**, the following is a theorem.

If 
$$A \subseteq B$$
 and  $B \subseteq A$  then  $A = B$ .

PROOF: If every element of A is an element of B, and every element of B is an element of A, then A and B have exactly the same elements.  $\Box$ 

**Proposition Schema 1.1.7.** For any classes A, B and C, the following is a theorem.

If 
$$A \subseteq B$$
 and  $B \subseteq C$  then  $A \subseteq C$ .

PROOF: If every element of A is an element of B and every element of B is an element of C then every element of A is an element of C.

### 1.1.2 Constructions of Classes

**Definition 1.1.8** (Empty Class). The *empty class*  $\emptyset$  is  $\{x \mid \bot\}$ . Every other class is *nonempty*.

**Definition 1.1.9** (Universal Class). The universal class V is  $\{x \mid \top\}$ .

**Definition 1.1.10** (Enumeration). Given objects  $a_1, \ldots, a_n$ , we define the class  $\{a_1, \ldots, a_n\}$  to be the class  $\{x \mid x = a_1 \lor \cdots \lor x = a_n\}$ .

**Definition 1.1.11** (Intersection). For any classes **A** and **B**, the *intersection*  $\mathbf{A} \cap \mathbf{B}$  is  $\{x \mid x \in \mathbf{A} \land x \in \mathbf{B}\}.$ 

**Definition 1.1.12** (Union). For any classes **A** and **B**, the *union*  $\mathbf{A} \cup \mathbf{B}$  is  $\{x \mid x \in \mathbf{A} \lor x \in \mathbf{B}\}.$ 

**Definition 1.1.13** (Relative Complement). Let **A** and **B** be classes. The *relative complement* of **B** in **A** is the class  $\mathbf{A} - \mathbf{B} := \{x \in \mathbf{A} \mid x \notin \mathbf{B}\}.$ 

**Definition 1.1.14** (Symmetric Difference). For any classes **A** and **B**, the *symmetric difference* is the class  $\mathbf{A} + \mathbf{B} := (\mathbf{A} - \mathbf{B}) \cup (\mathbf{B} - \mathbf{A})$ .

**Definition 1.1.15** (Pairwise disjoint). Let **A** be a class. We say the elements of **A** are *pairwise disjoint* iff, for all  $x, y \in \mathbf{A}$ , if  $x \cap y \neq \emptyset$  then x = y.

## 1.2 Sets and the Axiom of Extensionality

**Definition 1.2.1** (Axiom of Extensionality). The *Axiom of Extensionality* is the statement: if two sets have exactly the same members, then they are equal.

$$\forall x, y (\forall z (z \in x \Leftrightarrow z \in y) \Rightarrow x = y)$$
.

When working in a theory with the Axiom of Extensionality, we may identify a set a with the class  $\{x \mid x \in a\}$ . Our use of the symbols  $\in$  and = is consistent. We say a class  $\mathbf{A}$  is a set iff there exists a set a such that  $a = \mathbf{A}$ ; that is,  $\{x \mid \phi[x]\}$  is a set iff  $\exists a \forall x (x \in a \leftrightarrow \phi[x])$ . Otherwise,  $\mathbf{A}$  is a proper class.

**Definition 1.2.2** (Subset). If A is a set and  $A \subseteq \mathbf{B}$ , we say A is a *subset* of **B**.

**Definition 1.2.3** (Union). The *union* of a class **A** is  $\{x \mid \exists X \in \mathbf{A}.x \in X\}$ . We write  $\bigcup_{P(x)} t(x)$  for  $\bigcup \{t(x) \mid P(x)\}$ .

**Definition 1.2.4** (Intersection). The *intersection* of a class **A** is  $\{x \mid \forall X \in \mathbf{A}.x \in X\}$ . We write  $\bigcap_{P(x)} t(x)$  for  $\bigcap \{t(x) \mid P(x)\}$ .

**Definition 1.2.5** (Power Class). For any class **A**, the *power class*  $\mathcal{P}$ **A** is  $\{X \mid X \subseteq \mathbf{A}\}$ .

## 1.3 The Other Axioms

**Definition 1.3.1** (Empty Set Axiom). The *Empty Set Axiom* is the statement: The empty class  $\emptyset$  is a set.

$$\exists e \forall xx \notin e$$

**Definition 1.3.2** (Pairing Axiom). The *Pairing Axiom* or *Pair Set Axiom* is the statement: for any sets a and b, the class  $\{a, b\}$  is a set.

$$\forall a \forall b \exists c \forall x (x \in c \Leftrightarrow x = a \lor x = b)$$

**Definition 1.3.3** (Union Axiom). The *Union Axiom* is the statement: for any set A, the class  $\bigcup A$  is a set.

$$\forall A \exists B \forall x (x \in B \Leftrightarrow \exists y (y \in A \land x \in y))$$

**Definition 1.3.4** (Comprehension Axiom Scheme). The Comprehension Axiom Scheme, Subset Axiom Scheme or Aussonderungsaxiom Scheme is the set of sentences of the form, for any class **A**: If **A** is a subclass of a set then **A** is a set

That is, for any property  $P[x, y_1, \ldots, y_n]$ :

For any sets  $a_1, \ldots, a_n$  and B, the class  $\{x \in B \mid P[x, a_1, \ldots, a_n]\}$  is a set.

$$\forall a_1, \dots, a_n, B. \exists C. \forall x (x \in C \Leftrightarrow x \in B \land P[x, a_1, \dots, a_n])$$

**Definition 1.3.5** (Replacement Axiom Scheme). The Replacement Axiom Scheme is the set of sentences of the form, for some property  $P[x, y, z_1, \ldots, z_n]$ :

For any sets  $a_1, \ldots, a_n, B$ , assume for all  $x \in B$  there exists at most one y such that  $P[x, y, a_1, \ldots, a_n]$ . Then  $\{y \mid \exists x \in B. P[x, y, a_1, \ldots, a_n] \text{ is a set. }$ 

$$\forall a_1, \dots, a_n, B(\forall x \in B. \forall y, y'(P[x, y, a_1, \dots, a_n] \land P[x, y', a_1, \dots, a_n] \Rightarrow y = y') \Rightarrow$$
$$\exists C \forall y (y \in C \Leftrightarrow \exists x \in B. P[x, y, a_1, \dots, a_n]))$$

**Definition 1.3.6** (Power Set Axiom). The *Power Set Axiom* is the statement: the power class of a set is a set.

$$\forall A \exists B \forall x (x \in B \Leftrightarrow \forall y (y \in x \Rightarrow y \in A))$$

**Definition 1.3.7** (Axiom of Infinity). The *Axiom of Infinity* is the statement: there exists a set I such that  $\emptyset \in I$  and  $\forall x \in I.x \cup \{x\} \in I$ .

$$\exists I (\exists e \in I. \forall x. x \notin e \land \forall x \in I. \exists y \in I. \forall z (z \in y \Leftrightarrow z \in x \lor z = x))$$

**Definition 1.3.8** (Axiom of Choice). The *Axiom of Choice* is the statement: For any set A of pairwise disjoint, nonempty sets, there exists a set C such that, for all  $x \in A$ , we have  $x \cap C$  has exactly one element.

$$\forall A(\forall x \in A. \exists yy \in x \land \\ \forall x, y \in A. \forall z(z \in x \land z \in y \Rightarrow x = y) \Rightarrow \\ \exists C. \forall x \in A. \exists y \forall z(z \in x \land z \in C \Leftrightarrow z = y))$$

**Definition 1.3.9** (Axiom of Regularity). The Axiom of Regularity or Axiom of Foundation is the statement: for any A, if A has a member, then there exists  $m \in A$  such that  $m \cap A = \emptyset$ .

$$\forall A(\exists x. x \in A \Rightarrow \exists m \in A. \neg \exists x (x \in m \land x \in A))$$

**Definition 1.3.10** (Skolem Set Theory). Skolem set theory (S) is the theory whose axioms are:

- Extensionality
- Empty Set
- Pairing
- Union
- Comprehension
- Power Set

• Regularity

Let SC be the extension of S with the Axiom of Choice.

**Definition 1.3.11** (Skolem-Fraenkel Set Theory). Skolem-Fraenkel set theory (SF) is the theory whose axioms are:

- Extensionality
- Empty Set
- Union
- Replacement
- Power Set
- Regularity

Let SFC be the extension of SF with the Axiom of Choice.

**Definition 1.3.12** (Zermelo Set Theory). Zermelo set theory is the theory whose axioms are:

- Extensionality
- Pairing
- Union
- Comprehension
- Power Set
- Infinity
- Regularity

We label theorems with Z when they are provable in Zermelo set theory. Let ZC be the extension of Z with the Axiom of Choice.

**Definition 1.3.13** (Fraenkel-Mostowski Set Theory). Fraenkel-Mostowski set theory (FM) is the theory whose axioms are:

- The Axiom of Extensionality with Urelements: For any sets x and y, if x is nonempty and x and y have exactly the same elements, then x = y.
- Union
- Replacement
- Power Set
- Infinity

• Regularity

We write FMC for the extension of FM with Choice.

**Definition 1.3.14** (Zermelo-Fraenkel Set Theory). Zermelo-Fraenkel set theory is the theory whose axioms are:

- Extensionality
- Union
- Replacement
- Power Set
- Infinity
- Regularity

We label theorems with ZF when they are provable in Zermelo-Fraenkel set theory.

Let ZFC be the extension of ZF with the Axiom of Choice.

We label a theorem with FOL if it can be proved in first-order logic, i.e. from no axioms.

## 1.4 ZFC Extends Z

**Proposition 1.4.1.** The Axiom of Infinity implies the Empty Set Axiom.

Proof: Trivial.

Corollary 1.4.1.1. Z is an extension of S.

ZF is an extension of SF.

ZC is an extension of SC.

ZFC is an extension of SFC.

**Proposition 1.4.2.** The Axiom of Pairing is a theorem of SF.

#### Proof:

- $\langle 1 \rangle 1$ . Let: a, b be sets.
- $\langle 1 \rangle 2$ . Let: P(x,y) be the predicate  $(x = \emptyset \land y = a) \lor (x = \mathcal{P}\emptyset \land y = b)$ .
- $\langle 1 \rangle 3$ . For all  $x \in \mathcal{PP}\emptyset$ , there exists at most one y such that P(x,y).
  - $\langle 2 \rangle 1$ . Let:  $x \in \mathcal{PP}\emptyset$
  - $\langle 2 \rangle 2$ . Let: y and y' be sets.
  - $\langle 2 \rangle 3$ . Assume: P(x,y) and P(x,y')
  - $\langle 2 \rangle 4$ .  $(x = \emptyset \land y = a) \lor (x = \mathcal{P}\emptyset \land y = b)$

PROOF: From  $\langle 2 \rangle 3$ .

 $\langle 2 \rangle 5. \ (x = \emptyset \land y' = a) \lor (x = \mathcal{P}\emptyset \land y' = b)$ 

Proof: From  $\langle 2 \rangle 3$ .

```
 \begin{array}{l} \langle 2 \rangle 6. \; \emptyset \neq \mathcal{P} \emptyset \\ \quad \text{PROOF: Since } \emptyset \in \mathcal{P} \emptyset \text{ and } \emptyset \notin \emptyset. \\ \langle 2 \rangle 7. \; y = y' \\ \langle 1 \rangle 4. \; \text{Let: } A \text{ be the set } \{y \mid \exists x \in \mathcal{PP} \emptyset. P(x,y)\}. \\ \langle 1 \rangle 5. \; A = \{a,b\} \\ \sqcap \end{array}
```

**Proposition Schema 1.4.3.** Every instance of the Comprehension Axiom Scheme is a theorem of SF.

```
Proof:
\langle 1 \rangle 1. Let: P(x) be a predicate.
\langle 1 \rangle 2. Let: A be a set.
\langle 1 \rangle 3. Let: Q(x,y) be the predicate P(x) \wedge y = x.
\langle 1 \rangle 4. For all x \in A, there exists at most one y such that Q(x,y).
    \langle 2 \rangle 1. Let: x \in A
    \langle 2 \rangle 2. Let: y and y' be sets.
    \langle 2 \rangle 3. Assume: Q(x,y) and Q(x,y')
   \langle 2 \rangle 4. \ x \in A \land P(x) \land y = x \land y' = x
       PROOF: From \langle 2 \rangle 3.
    \langle 2 \rangle 5. \ y = y'
       PROOF: From \langle 2 \rangle 4.
\langle 1 \rangle5. Let: B be the set \{y \mid \exists x \in A.Q(x,y)\}
   PROOF: This is a set by an Axiom of Replacement and \langle 1 \rangle 4.
\langle 1 \rangle 6. \ B = \{ y \in A \mid P(y) \}
   Proof:
                        y \in B \Leftrightarrow \exists x \in A.Q(x,y)
                                                                                              (\langle 1 \rangle 5)
                                  \Leftrightarrow \exists x \in A(P(x) \land y = x)
                                                                                              (\langle 1 \rangle 3)
```

Corollary 1.4.3.1. SF is an extension of S.

 $\Leftrightarrow P(y)$ 

SFC is an extension of SC. ZF is an extension of Z. ZFC is an extension of ZC.

П

# 1.5 Consequences of the Axioms

**Proposition 1.5.1** (S). The union of two sets is a set.

PROOF: Because  $A \cup B = \bigcup \{A, B\}$ .  $\square$ 

**Proposition Schema 1.5.2** (S). For any number n, the following is a theorem: For any sets  $a_1, \ldots, a_n$ , the class  $\{a_1, \ldots, a_n\} = \{x \mid x = a_1 \lor \cdots \lor x = a_n\}$  is a set.

```
PROOF: The case n = 1 follows from Pairing since \{a\} = \{a, a\}.
If we have proved the theorem for n we have \{a_1, \ldots, a_n, a_{n+1}\} = \{a_1, \ldots, a_n\} \cup
\{a_{n+1}\}.
Proposition 1.5.3 (S). No set is a member of itself.
Proof:
\langle 1 \rangle 1. Let: x be any set.
\langle 1 \rangle 2. PICK m \in \{x\} such that m \cap \{x\} = \emptyset.
  PROOF: Axiom of Regularity.
\langle 1 \rangle 3. \ m = x
\langle 1 \rangle 4. \ x \cap \{x\} = \emptyset
\langle 1 \rangle 5. \ x \notin x
Corollary 1.5.3.1 (S). The universal class V is a proper class.
PROOF: If V is a set then V \in V, contradicting the Proposition. \square
Proposition 1.5.4 (S). There are no sets a and b such that a \in b and b \in a.
Proof:
\langle 1 \rangle 1. Let: a and b be any sets.
\langle 1 \rangle 2. PICK m \in \{a, b\} such that m \cap \{a, b\} = \emptyset
\langle 1 \rangle 3. Case: m=a
  PROOF: Then b \notin a.
\langle 1 \rangle 4. Case: m = b
  PROOF: Then a \notin b.
Proposition 1.5.5 (S). The intersection of a set and a class is a set.
PROOF: Immediate from Comprehension.
Proposition 1.5.6 (S). The relative complement of a class in a set is a set.
Proof: Immediate from Comprehension.
Corollary 1.5.6.1 (S). The symmetric difference of two sets is a set.
Proposition 1.5.7 (S). The intersection of a nonempty class is a set.
Proof:
\langle 1 \rangle 1. Let: A be a nonempty class.
\langle 1 \rangle 2. Pick B \in \mathbf{A}
\langle 1 \rangle 3. \cap \mathbf{A} \subseteq B
\langle 1 \rangle 4. \bigcap \mathbf{A} is a set.
  PROOF: By Comprehension.
```

**Proposition Schema 1.5.8** (FOL). For any classes **A** and **B**, the following is a theorem:

If 
$$A \subseteq B$$
 then  $\mathcal{P}A \subseteq \mathcal{P}B$ .

PROOF: Every subset of **A** is a subset of **B**.  $\square$ 

**Proposition Schema 1.5.9** (FOL). For any classes **A** and **B**, the following is a theorem:

If 
$$\mathbf{A} \subseteq \mathbf{B}$$
 then  $\bigcup \mathbf{A} \subseteq \bigcup \mathbf{B}$ .

PROOF: If  $x \in X \in \mathbf{A}$  then  $x \in X \in \mathbf{B}$ .  $\square$ 

**Proposition Schema 1.5.10** (S). For any class **A**, the following is a theorem:

$$\mathbf{A} = \bigcup \mathcal{P} \mathbf{A}$$

Proof:

 $\langle 1 \rangle 1$ .  $\mathbf{A} \subseteq \bigcup \mathcal{P} \mathbf{A}$ 

PROOF: For all  $x \in \mathbf{A}$  we have  $x \in \{x\} \in \mathcal{P}\mathbf{A}$ .

 $\langle 1 \rangle 2$ .  $\bigcup \mathcal{P} \mathbf{A} \subseteq \mathbf{A}$ 

- $\langle 2 \rangle 1$ . Let:  $x \in \bigcup \mathcal{P} \mathbf{A}$
- $\langle 2 \rangle 2$ . PICK  $X \in \mathcal{P}\mathbf{A}$  such that  $x \in X$
- $\langle 2 \rangle 3. \ X \subseteq \mathbf{A}$
- $\langle 2 \rangle 4. \ x \in \mathbf{A}$

## 1.6 Transitive Classes

**Definition 1.6.1** (Transitive Class). A class **A** is a *transitive class* iff whenever  $x \in y \in \mathbf{A}$  then  $x \in \mathbf{A}$ .

**Proposition Schema 1.6.2** (FOL). For any class **A**, the following is a theorem:

The following are equivalent.

- 1. A is a transitive class.
- $2. \mid JA \subseteq A$
- 3. Every element of A is a subset of A.
- 4.  $\mathbf{A} \subseteq \mathcal{P}\mathbf{A}$

PROOF: Immediate from definitions.  $\Box$ 

**Proposition Schema 1.6.3** (FOL). For any class **A**, the following is a theorem:

If **A** is a transitive class then  $\bigcup \mathbf{A}$  is a transitive class.

Proof:

```
\langle 1 \rangle 1. Assume: A is a transitive class.
\langle 1 \rangle 2. Let: x \in y \in \bigcup \mathbf{A}
\langle 1 \rangle 3. \ y \in \mathbf{A}
   PROOF: Since \bigcup \mathbf{A} \subseteq \mathbf{A} by Proposition 1.6.2.
\langle 1 \rangle 4. \ x \in \bigcup \mathbf{A}
Proposition Schema 1.6.4 (S). For any class A, the following is a theorem:
      We have A is a transitive class if and only if \mathcal{P}\mathbf{A} is a transitive class.
Proof:
\langle 1 \rangle 1. If A is a transitive class then \mathcal{P}A is a transitive class.
    \langle 2 \rangle 1. Assume: A is a transitive class.
    \langle 2 \rangle 2. \mathbf{A} \subseteq \mathcal{P} \mathbf{A}
       Proof: Proposition 1.6.2.
    \langle 2 \rangle 3. \mathcal{P} \mathbf{A} \subseteq \mathcal{P} \mathcal{P} \mathbf{A}
       Proof: Proposition 1.5.8.
    \langle 2 \rangle 4. \mathcal{P}\mathbf{A} is a transitive class.
       Proof: Proposition 1.6.2.
\langle 1 \rangle 2. If \mathcal{P}\mathbf{A} is a transitive class then \mathbf{A} is a transitive class.
    \langle 2 \rangle 1. Assume: \mathcal{P}\mathbf{A} is a transitive class.
    \langle 2 \rangle 2. \bigcup \mathcal{P} \mathbf{A} \subseteq \mathcal{P} \mathbf{A}
       Proof: Proposition 1.6.2.
    \langle 2 \rangle 3. \mathbf{A} \subseteq \mathcal{P} \mathbf{A}
       Proof: Proposition 1.5.10.
    \langle 2 \rangle 4. A is a transitive class.
       Proof: Proposition 1.6.2.
Proposition Schema 1.6.5 (FOL). For any class A, the following is a theo-
     If every member of A is a transitive set then \bigcup A is a transitive class.
Proof:
\langle 1 \rangle 1. Assume: Every member of A is a transitive set.
\langle 1 \rangle 2. Let: x \in y \in \bigcup \mathbf{A}
\langle 1 \rangle 3. Pick A \in \mathbf{A} such that y \in A.
\langle 1 \rangle 4. \ x \in A
   PROOF: Since A is a transitive set.
\langle 1 \rangle 5. \ x \in \bigcup \mathbf{A}
Proposition Schema 1.6.6 (FOL). For any class A, the following is a theo-
```

If every member of **A** is a transitive set then  $\bigcap \mathbf{A}$  is a transitive class.

### Proof:

 $\langle 1 \rangle 1$ . Assume: Every member of **A** is a transitive set.

## 1.6. TRANSITIVE CLASSES

 $\begin{array}{l} \langle 1 \rangle 2. \ \ \mathrm{Let:} \ \ x \in y \in \bigcap \mathbf{A} \\ \ \ \ \mathrm{Prove:} \ \ \ x \in \bigcap \mathbf{A} \\ \langle 1 \rangle 3. \ \ \mathrm{Let:} \ \ A \in \mathbf{A} \end{array}$ 

 $\langle 1 \rangle 4. \ y \in A$  $\langle 1 \rangle 5. \ x \in A$ 

PROOF: Since A is a transitive set.

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# Chapter 2

# Relations

## 2.1 Ordered Pairs

**Definition 2.1.1** (Ordered Pair). For any sets a and b, the *ordered pair* (a, b) is defined to be  $\{\{a\}, \{a, b\}\}.$ 

**Theorem 2.1.2** (S). For any sets a, b, c, d, we have (a,b) = (c,d) if and only if a = c and b = d.

#### Proof:

```
\langle 1 \rangle 1. If (a, b) = (c, d) then a = c and b = d.
    \langle 2 \rangle 1. Assume: \{\{a\}, \{a, b\}\} = \{\{c\}, \{c, d\}\}
    \langle 2 \rangle 2. \cap \{\{a\}, \{a, b\}\} = \bigcap \{\{c\}, \{c, d\}\}\
    \langle 2 \rangle 3. \{a\} = \{c\}
    \langle 2 \rangle 4. a = c
    \langle 2 \rangle 5. \bigcup \{\{a\}, \{a, b\}\} = \bigcup \{\{c\}, \{c, d\}\}\
    \langle 2 \rangle 6. \ \{a,b\} = \{c,d\}
    \langle 2 \rangle 7. b = c or b = d
    \langle 2 \rangle 8. a = d or b = d
    \langle 2 \rangle 9. If b = c and a = d then b = d
        Proof: By \langle 2 \rangle 4.
    \langle 2 \rangle 10. b=d
        PROOF: From \langle 2 \rangle 7, \langle 2 \rangle 8, \langle 2 \rangle 9.
\langle 1 \rangle 2. If a = c and b = d then (a, b) = (c, d).
    PROOF: First-order logic.
```

**Definition 2.1.3** (Cartesian Product). The *Cartesian product* of classes **A** and **B** is the class  $\mathbf{A} \times \mathbf{B} := \{(x,y) \mid x \in \mathbf{A}, y \in \mathbf{B}\}.$ 

**Proposition 2.1.4** (S). For any sets A and B, the class  $A \times B$  is a set.

PROOF: It is a subset of  $\mathcal{PP}(A \cup B)$ .  $\square$ 

**Proposition Schema 2.1.5** (S). For any classes **A**, **B** and **C**, the following is a theorem:

$$\mathbf{A} \times (\mathbf{B} \cup \mathbf{C}) = (\mathbf{A} \times \mathbf{B}) \cup (\mathbf{A} \times \mathbf{C})$$

Proof:

$$(x,y) \in \mathbf{A} \times (\mathbf{B} \cup \mathbf{C}) \Leftrightarrow x \in \mathbf{A} \wedge (y \in \mathbf{B} \vee y \in \mathbf{C})$$
$$\Leftrightarrow (x \in \mathbf{A} \wedge y \in \mathbf{B}) \vee (x \in \mathbf{A} \wedge y \in \mathbf{C})$$
$$\Leftrightarrow (x,y) \in (\mathbf{A} \times \mathbf{B}) \cup (\mathbf{A} \times \mathbf{C})$$

**Proposition Schema 2.1.6** (S). For any classes **A** and **B**, the following is a theorem:

If 
$$\mathbf{A} \times \mathbf{B} = \mathbf{A} \times \mathbf{C}$$
 and  $\mathbf{A}$  is nonempty then  $\mathbf{B} = \mathbf{C}$ .

Proof:

- $\langle 1 \rangle 1$ . Pick  $a \in \mathbf{A}$
- $\langle 1 \rangle 2$ . For all x we have  $x \in \mathbf{B}$  iff  $x \in \mathbf{C}$ .

Proof:

$$x \in \mathbf{B} \Leftrightarrow (a, x) \in \mathbf{A} \times \mathbf{B}$$
  
 $\Leftrightarrow (a, x) \in \mathbf{A} \times \mathbf{C}$   
 $\Leftrightarrow x \in \mathbf{C}$ 

**Proposition Schema 2.1.7** (S). For any classes **A** and **B**, the following is a theorem:

$$\mathbf{A} \times \bigcup \mathbf{B} = \{(a, b) \mid \exists Y \in \mathbf{B}. (a \in \mathbf{A} \land b \in Y)\}\$$

Proof:

$$(x,y) \in A \times \bigcup \mathbf{B} \Leftrightarrow x \in A \land \exists Y \in \mathbf{B}.y \in Y$$
$$\Leftrightarrow \exists Y \in \mathbf{B}(x \in A \land y \in Y)$$

## 2.2 Relations

**Definition 2.2.1** (Relation). A relation is a class of ordered pairs.

A relation **R** between classes **A** and **B** is a subclass of  $\mathbf{A} \times \mathbf{B}$ .

A (binary) relation on  $\mathbf{A}$  is a relation between  $\mathbf{A}$  and  $\mathbf{A}$ .

We write  $x\mathbf{R}y$  for  $(x,y) \in \mathbf{R}$ .

**Definition 2.2.2** (Domain). The *domain* of a class  $\mathbf{R}$  is the class

$$\operatorname{dom} \mathbf{R} := \{ x \mid \exists y. x \mathbf{R} y \} .$$

**Proposition 2.2.3** (S). The domain of a set is a set.

**Definition 2.2.4** (Range). The range of a class  $\mathbf{R}$  is the class

$$\operatorname{ran} \mathbf{R} := \{ y \mid \exists x. x \mathbf{R} y \}$$

**Proposition 2.2.5** (S). The range of a set is a set.

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## 2.2.1 Identity Functions

**Definition 2.2.6** (Identity Function). For any class A, the *identity function* or *diagonal relation*  $id_A$  on A is

$$id_{\mathbf{A}} := \{(x, x) \mid x \in \mathbf{A}\} .$$

## 2.2.2 Inverses

**Definition 2.2.7** (Inverse). The *inverse* of a relation  $\mathbf{R}$  between  $\mathbf{A}$  and  $\mathbf{B}$  is the relation  $\mathbf{R}^{-1}$  between  $\mathbf{B}$  and  $\mathbf{A}$  defined by

$$b\mathbf{R}^{-1}a \Leftrightarrow a\mathbf{R}b$$
.

**Proposition Schema 2.2.8** (S). For any classes A, B and R, the following is a theorem:

If **R** is a relation between **A** and **B** then  $(\mathbf{R}^{-1})^{-1} = \mathbf{R}$ .

Proof:

$$x(\mathbf{R}^{-1})^{-1}y \Leftrightarrow y\mathbf{R}^{-1}x$$
  
 $\Leftrightarrow x\mathbf{R}y$ 

## 2.2.3 Composition

**Definition 2.2.9** (Composition). Let  $\mathbf{R}$  be a relation between  $\mathbf{A}$  and  $\mathbf{B}$ , and  $\mathbf{S}$  be a relation between  $\mathbf{B}$  and  $\mathbf{C}$ . The *composition*  $\mathbf{S} \circ \mathbf{R}$  is the relation between  $\mathbf{A}$  and  $\mathbf{C}$  defined by

$$a(\mathbf{S} \circ \mathbf{R})c \Leftrightarrow \exists b(a\mathbf{R}b \wedge b\mathbf{S}c)$$
.

**Proposition Schema 2.2.10** (S). For any classes A, B, C, R and S, the following is a theorem:

If  ${\bf R}$  is a relation between  ${\bf A}$  and  ${\bf B}$ , and  ${\bf S}$  is a relation between  ${\bf B}$  and  ${\bf C}$ , then

$$(\mathbf{S} \circ \mathbf{R})^{-1} = \mathbf{R}^{-1} \circ \mathbf{S}^{-1} .$$

Proof:

$$z(\mathbf{S} \circ \mathbf{R})^{-1}x \Leftrightarrow x(\mathbf{S} \circ \mathbf{R})z$$

$$\Leftrightarrow \exists y.(x\mathbf{R}y \wedge y\mathbf{S}z)$$

$$\Leftrightarrow \exists y.(y\mathbf{R}^{-1}x \wedge z\mathbf{S}^{-1}y)$$

$$\Leftrightarrow z(\mathbf{R}^{-1} \circ \mathbf{S}^{-1})x$$

## 2.2.4 Properties of Relaitons

**Definition 2.2.11** (Reflexive). Let **R** be a binary relation on **A**. Then **R** is reflexive on **A** iff  $\forall x \in \mathbf{A}.(x,x) \in \mathbf{R}$ .

**Proposition Schema 2.2.12** (S). For any classes A and R, the following is a theorem:

If **R** is a reflexive relation on **A** then so is  $\mathbf{R}^{-1}$ .

### Proof:

- $\langle 1 \rangle 1$ . Let:  $x \in \mathbf{A}$
- $\langle 1 \rangle 2$ .  $x \mathbf{R} x$

PROOF: Since  $\mathbf{R}$  is reflexive.

 $\langle 1 \rangle 3. \ x \mathbf{R}^{-1} x$ 

**Definition 2.2.13** (Irreflexive). A relation **R** is *irreflexive* iff there is no x such that  $(x, x) \in \mathbf{R}$ .

**Definition 2.2.14** (Symmetric). A relation **R** is *symmetric* iff, whenever  $x\mathbf{R}y$ , then  $y\mathbf{R}x$ .

**Definition 2.2.15** (Antisymmetric). A relation **R** is *antisymmetric* iff, whenever  $x\mathbf{R}y$  and  $y\mathbf{R}x$ , then x=y.

**Proposition Schema 2.2.16** (S). For any classes  ${\bf A}$  and  ${\bf R}$ , the following is a theorem:

If **R** is an antisymmetric relation on **A** then so is  $\mathbf{R}^{-1}$ .

#### PROOF

- $\langle 1 \rangle 1$ . Assume:  $x \mathbf{R}^{-1} y$  and  $y \mathbf{R}^{-1} x$
- $\langle 1 \rangle 2$ .  $y \mathbf{R} x$  and  $x \mathbf{R} y$
- $\langle 1 \rangle 3. \ x = y$

PROOF: Since  $\mathbf{R}$  is antisymmetric.

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**Definition 2.2.17** (Transitive). A relation **R** is *transitive* iff, whenever  $x\mathbf{R}y$  and  $y\mathbf{R}z$ , then  $x\mathbf{R}z$ .

**Proposition Schema 2.2.18** (S). For any classes **A**, **B** and **R**, the following is a theorem:

If **R** is a transitive relation between **A** and **B** then  $\mathbf{R}^{-1}$  is transitive.

## Proof:

- $\langle 1 \rangle 1$ . Assume:  $(x,y), (y,z) \in \mathbf{R}^{-1}$
- $\langle 1 \rangle 2. \ (y, x), (z, y) \in \mathbf{R}$
- $\langle 1 \rangle 3. \ (z,x) \in \mathbf{R}$
- $\langle 1 \rangle 4. \ (x,z) \in \mathbf{R}^{-1}$

**Proposition 2.2.19** (S). For any relation R on a set A, there exists a smallest transitive relation on A that includes R.

PROOF: The relation is  $\bigcap \{S \in \mathcal{P}A^2 \mid R \subseteq S, S \text{ is transitive}\}$ .  $\square$ 

**Definition 2.2.20** (Transitive Closure). For any relation R on a set A, the transitive closure of R is the smallest transitive relation that includes R.

**Definition 2.2.21** (Minimal). Let **R** be a relation on **A**. An element  $m \in \mathbf{A}$  is *minimal* iff there is no  $x \in \mathbf{A}$  such that  $x\mathbf{R}m$ .

**Definition 2.2.22** (Maximal). Let **R** be a relation on **A**. An element  $m \in \mathbf{A}$  is *maximal* iff there is no  $x \in \mathbf{A}$  such that  $m\mathbf{R}x$ .

# 2.3 n-ary Relations

**Definition Schema 2.3.1.** For any sets  $a_1, \ldots, a_n$ , define the *ordered n-tuple*  $(a_1, \ldots, a_n)$  by

$$(a_1) := a_1$$
  
 $(a_1, \dots, a_n, a_{n+1}) = ((a_1, \dots, a_n), a_{n+1})$ 

**Definition Schema 2.3.2.** An n-ary relation on A is a class of ordered n-tuples all of whose components are in A.

## 2.4 Well Founded Relations

**Definition 2.4.1** (Well Founded). A relation **R** is well founded iff:

- Every nonempty set has an **R**-minimal element.
- For every set x, there exists a set u such that  $x \subseteq u$  and, for all w, y, if  $y \in u$  and  $w\mathbf{R}y$  then  $w \in u$ .

**Proposition Schema 2.4.2** (S). For any class  $\mathbf{R}$ , the following is a theorem: Assume  $\mathbf{R}$  is a well founded relation. For any set a, the class  $\{x \mid x\mathbf{R}a\}$  is a set.

### Proof:

 $\langle 1 \rangle 1$ . PICK a set u such that  $\{a\} \subseteq u$  and, for all w, y, if  $y \in u$  and  $w \mathbf{R} y$  then  $w \in u$ .

$$w \in u.$$

$$\langle 1 \rangle 2. \{x \mid x\mathbf{R}a\} \subseteq u$$

**Proposition Schema 2.4.3** (S). For any classes **A**, **B** and **R**, the following is a theorem:

Assume **R** is a well founded relation on **A** and **B**  $\subseteq$  **A** is nonempty. Then **B** has an **R**-minimal element.

### Proof:

- $\langle 1 \rangle 1$ . Pick  $b \in \mathbf{B}$
- $\langle 1 \rangle 2$ . Let:  $S = \{ x \in \mathbf{B} \mid x\mathbf{R}b \}$
- $\langle 1 \rangle 3$ . Case:  $S = \emptyset$

PROOF: In this case b is an **R**-minimal element of **B**.

 $\langle 1 \rangle 4$ . Case:  $S \neq \emptyset$ 

PROOF: In this cases S has an  ${\bf R}$ -minimal element, which is an  ${\bf R}$ -minimal element of  ${\bf B}$ .

**Proposition Schema 2.4.4** (FOL). For any classes  $\mathbf{R}$  and  $\mathbf{S}$ , the following is a theorem:

Assume **R** is a relation. If **S** is a well-founded relation and  $\mathbf{R} \subseteq \mathbf{S}$  then **R** is a well founded relation.

PROOF: Immediate from definitions.  $\Box$ 

**Theorem Schema 2.4.5** (Transfinite Induction Principle (S)). For any classes **A**, **B** and **R**, the following is a theorem:

Assume **R** is a well founded relation on **A** and **B**  $\subseteq$  **A**. Assume that, for all  $t \in$  **A**,

$$\{x \in \mathbf{A} \mid x\mathbf{R}t\} \subseteq \mathbf{B} \Rightarrow t \in \mathbf{B}$$
.

Then  $\mathbf{B} = \mathbf{A}$ .

#### Proof:

- $\langle 1 \rangle 1$ . Assume: for a contradiction  $\mathbf{B} \neq \mathbf{A}$
- $\langle 1 \rangle 2$ . PICK an **R**-minimal element m of  $\mathbf{A} \mathbf{B}$ .

Proof: Proposition 2.4.3.

 $\langle 1 \rangle 3. \{ x \in \mathbf{A} \mid x\mathbf{R}m \} \subseteq \mathbf{B}$ 

PROOF: By minimality of m.

- $\langle 1 \rangle 4. \ m \in \mathbf{B}$
- $\langle 1 \rangle 5$ . Q.E.D.

Proof: This is a contradiction.

**Theorem 2.4.6** (S). The transitive closure of a well founded relation on a set is well founded.

### Proof:

- $\langle 1 \rangle 1$ . Let: R be a well founded relation on the set A.
- $\langle 1 \rangle 2$ . Let:  $R^t$  be the transitive closure of R.
- $\langle 1 \rangle 3$ . For any  $x, y \in A$ , if  $xR^ty$  then there exists  $z \in A$  such that zRy. PROOF:  $\{(x,y) \in A^2 \mid \exists z \in A. zRy\}$  is a transitive relation on A that includes R.
- $\langle 1 \rangle 4$ . Let: B be a nonempty subset of A.
- $\langle 1 \rangle$ 5. Pick an *R*-minimal element *b* of *B*.
- $\langle 1 \rangle 6$ . b is  $R^t$ -minimal in B.

PROOF: If there exists x such that  $xR^tb$  then there exists z such that zRb by  $\langle 1 \rangle 3$ .

**Definition 2.4.7** (Initial Segment). Let **R** be a relation on **A** and  $a \in \mathbf{A}$ . The *initial segment* up to a is

$$\operatorname{seg} a := \{ x \in \mathbf{A} \mid x\mathbf{R}a \} .$$

**Theorem Schema 2.4.8** (Transfinite Recursion Theorem Schema (ZFC)). For any classes  $\mathbf{A}$ ,  $\mathbf{R}$  and any property G[x, y, z], there exists a class  $\mathbf{F}$  such that, for any class  $\mathbf{F}'$  the following is a theorem:

Assume that **R** is a well-founded relation on **A**. Assume that, for any f and t, there exists a unique z such that G[f,t,z]. Then  $\mathbf{F}: \mathbf{A} \to \mathbf{V}$  such that, for all  $t \in \mathbf{A}$ , we have  $\mathbf{F} \upharpoonright \operatorname{seg} t$  is a set and

$$G[\mathbf{F} \upharpoonright \operatorname{seg} t, t, \mathbf{F}(t)]$$
.

If  $\mathbf{F}' : \mathbf{A} \to \mathbf{V}$  satisfies that, for all  $t \in \mathbf{A}$ , we have  $\mathbf{F}' \upharpoonright \operatorname{seg} t$  is a set and  $G[\mathbf{F}' \upharpoonright \operatorname{seg} t, t, \mathbf{F}'(t)]$ , then  $\mathbf{F}' = \mathbf{F}$ .

#### Proof:

- $\langle 1 \rangle 1$ . For B a subset of A, let us say a function  $v : B \to V$  is acceptable iff, for all  $x \in B$ , we have  $\operatorname{seg} x \subseteq B$  and  $G[v \upharpoonright \operatorname{seg} x, x, v(x)]$
- $\langle 1 \rangle 2$ . Let: **K** be the class of all acceptable functions.
- $\langle 1 \rangle 3$ . Let:  $\mathbf{F} = \bigcup \mathbf{K}$
- $\langle 1 \rangle 4$ . For all  $B, C \subseteq \mathbf{A}$ , given  $v_1 : B \to \mathbf{V}$  and  $v_2 : C \to \mathbf{V}$  acceptable and  $x \in B \cap C$ , we have  $v_1(x) = v_2(x)$ 
  - $\langle 2 \rangle 1$ . Assume: as transfinite induction hypothesis  $\forall y \mathbf{R} x. y \in B \cap C \Rightarrow v_1(y) = v_2(y)$
  - $\langle 2 \rangle 2$ .  $v_1 \upharpoonright \operatorname{seg} x = v_2 \upharpoonright \operatorname{seg} x$
  - $\langle 2 \rangle 3$ .  $G[v_1 \upharpoonright \operatorname{seg} x, x, v_1(x)]$
  - $\langle 2 \rangle 4$ .  $G[v_2 \upharpoonright \operatorname{seg} x, x, v_2(x)]$
  - $\langle 2 \rangle 5. \ v_1(x) = v_2(x)$
- $\langle 1 \rangle 5$ . **F** is a function.
  - $\langle 2 \rangle 1$ . Assume:  $(x,y), (x,z) \in \mathbf{F}$
  - $\langle 2 \rangle 2$ . PICK acceptable  $v_1 : B \to \mathbf{V}$  and  $v_2 : C \to \mathbf{V}$  such that  $v_1(x) = y$  and  $v_2(x) = z$
  - $\langle 2 \rangle 3. \ y=z$

Proof: By  $\langle 1 \rangle 4$ .

- $\langle 1 \rangle 6$ . For all  $t \in \text{dom } \mathbf{F}$ , we have  $\mathbf{F} \upharpoonright \text{seg } t$  is a set and  $G[\mathbf{F} \upharpoonright \text{seg } t, t, \mathbf{F}(t)]$ 
  - $\langle 2 \rangle 1$ . Let:  $t \in \text{dom } \mathbf{F}$
  - $\langle 2 \rangle 2$ . PICK an acceptable  $v: A \to \mathbf{V}$  such that  $t \in A$
  - $\langle 2 \rangle 3$ . For all  $y \mathbf{R} x$  we have  $v(y) = \mathbf{F}(y)$
  - $\langle 2 \rangle 4$ . **F**  $\upharpoonright \operatorname{seg} x = v \upharpoonright \operatorname{seg} x$
  - $\langle 2 \rangle 5$ .  $G[v \upharpoonright \operatorname{seg} x, x, v(x)]$
  - $\langle 2 \rangle 6. \ G[\mathbf{F} \upharpoonright \operatorname{seg} x, x, \mathbf{F}(x)]$
- $\langle 1 \rangle 7$ . dom  $\mathbf{F} = \mathbf{A}$ 
  - $\langle 2 \rangle 1$ . Let:  $x \in \mathbf{A}$
  - $\langle 2 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall y \mathbf{R} x. y \in \mathbf{A}$
  - $\langle 2 \rangle 3$ . Assume: for a contradiction  $x \notin \text{dom } \mathbf{F}$

```
\langle 2 \rangle 4. F \upharpoonright \operatorname{seg} x is a set
         PROOF: Axiom of Replacement.
     \langle 2 \rangle 5. F \upharpoonright \operatorname{seg} x is acceptable
     \langle 2 \rangle 6. Let: y be the unique object such that G[\mathbf{F} \upharpoonright \operatorname{seg} x, x, y]
     \langle 2 \rangle 7. F \upharpoonright \operatorname{seg} x \cup \{(x,y)\} is acceptable
     \langle 2 \rangle 8. \ x \in \text{dom } \mathbf{F}
     \langle 2 \rangle 9. Q.E.D.
         PROOF: This is a contradiction.
\langle 1 \rangle 8. If \mathbf{F}' : \mathbf{A} \to \mathbf{V} satisfies the theorem, then \mathbf{F}' = \mathbf{F}.
     \langle 2 \rangle 1. Let: x \in \mathbf{A}
                 PROVE: \mathbf{F}'(x) = \mathbf{F}(x)
     \langle 2 \rangle 2. Assume: as transfinite induction hypothesis \forall y \mathbf{R} x. \mathbf{F}'(y) = \mathbf{F}(y)
     \langle 2 \rangle 3. \mathbf{F} \upharpoonright x = \mathbf{F}' \upharpoonright x
     \langle 2 \rangle 4. G[\mathbf{F} \upharpoonright x, x, \mathbf{F}(x)]
    \langle 2 \rangle 5. G[\mathbf{F}' \upharpoonright x, x, \mathbf{F}'(x)]
    \langle 2 \rangle 6. \ \mathbf{F}(x) = \mathbf{F}'(x)
```

# Chapter 3

# **Functions**

## 3.1 Functions

**Definition 3.1.1** (Function). A function from **A** to **B** is a relation **F** between **A** and **B** such that, for all  $x \in \mathbf{A}$ , there is only one y such that  $x\mathbf{F}y$ . We denote this y by  $\mathbf{F}(x)$ .

A binary operation on a class **A** is a function  $\mathbf{A}^2 \to \mathbf{A}$ .

**Definition 3.1.2** (Closed). Let  $\mathbf{F} : \mathbf{A} \to \mathbf{A}$  be a function and  $\mathbf{B} \subseteq \mathbf{A}$ . Then  $\mathbf{B}$  is *closed* under  $\mathbf{F}$  iff  $\forall x \in \mathbf{B}.\mathbf{F}(x) \in \mathbf{B}$ .

**Proposition 3.1.3** (S). For any class **A**, the following is a theorem:

$$\mathrm{id}_A:A\to A$$

PROOF: For all  $x \in \mathbf{A}$ , the only y such that  $(x, y) \in \mathrm{id}_{\mathbf{A}}$  is y = x.  $\square$ 

**Proposition Schema 3.1.4** (S). For any classes A, B, C, F and G, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{G}: \mathbf{B} \to \mathbf{C}$ . Then  $\mathbf{G} \circ \mathbf{F}: \mathbf{A} \to \mathbf{C}$  and, for all  $x \in \mathbf{A}$ , we have

$$(\mathbf{G} \circ \mathbf{F})(x) = \mathbf{G}(\mathbf{F}(x))$$
.

Proof:

 $\langle 1 \rangle 1. \ \forall x \in \mathbf{A}.(x, \mathbf{G}(\mathbf{F}(x))) \in \mathbf{G} \circ \mathbf{F})$ PROOF: Because  $(x, \mathbf{F}(x)) \in \mathbf{F}$  and  $(\mathbf{F}(x), \mathbf{G}(\mathbf{F}(x))) \in \mathbf{G}$ .  $\langle 1 \rangle 2. \ \text{If } (x, z) \in \mathbf{F} \circ \mathbf{G} \text{ then } z = \mathbf{G}(\mathbf{F}(x))$ 

 $\langle 2 \rangle 1$ . PICK  $y \in \mathbf{B}$  such that  $x \mathbf{F} y$  and  $y \mathbf{G} z$ 

 $\langle 2 \rangle 2. \ y = \mathbf{F}(x)$ 

 $\langle 2 \rangle 3. \ z = \mathbf{G}(y)$ 

 $\langle 2 \rangle 4. \ z = \mathbf{G}(\mathbf{F}(x))$ 

**Proposition 3.1.5** (SC). For any set A there exists a function  $F : \mathcal{P}A - \{\emptyset\} \to A$  (a choice function for A) such that, for every nonempty  $B \subseteq A$ , we have  $F(B) \in B$ .

```
Proof:
 \langle 1 \rangle 1. Let: A be a set.
 \langle 1 \rangle 2. Let: \mathcal{A} = \{ \{B\} \times B \mid B \in \mathcal{P}A - \{\emptyset\} \}
 \langle 1 \rangle 3. Every member of \mathcal{A} is nonempty.
 \langle 1 \rangle 4. Any two distinct members of \mathcal{A} are disjoint.
 \langle 1 \rangle5. PICK a set C such that, for all X \in \mathcal{A}, we have C \cap X is a singleton.
    Proof: Axiom of Choice.
 \langle 1 \rangle 6. Let: F = C \cap \bigcup \mathcal{A}
 \langle 1 \rangle 7. \ F : \mathcal{P}A - \{\emptyset\} \to A
     \langle 2 \rangle 1. F is a function.
         (3)1. Let: (B, b), (B, b') \in F
         \langle 3 \rangle 2. \ (B,b), (B,b') \in \{B\} \times B
             PROOF: Since (B, b), (B, b') \in \bigcup A.
         (3)3. (B,b), (B,b') \in C \cap (\{B\} \times B)
         \langle 3 \rangle 4. \ (B,b) = (B,b')
             PROOF: From \langle 1 \rangle 5.
         \langle 3 \rangle 5. \ b = b'
     \langle 2 \rangle 2. dom F = \mathcal{P}A - \{\emptyset\}
        Proof:
                     B \in \operatorname{dom} F
                 \Leftrightarrow \exists b.(B,b) \in F
                 \Leftrightarrow \exists b. ((B,b) \in \bigcup \mathcal{A} \land (B,b) \in C)
                 \Leftrightarrow \exists b. \exists B' \in \mathcal{P}A - \{\emptyset\}. ((B,b) \in \{B'\} \times B' \land (B,b) \in C)
                 \Leftrightarrow B \in \mathcal{P}A - \{\emptyset\} \land \exists b \in B.(B,b) \in C
                 \Leftrightarrow B \in \mathcal{P}A - \{\emptyset\}
                                                                                                                            (\langle 1 \rangle 5)
     \langle 2 \rangle 3. ran F \subseteq A
\langle 1 \rangle 8. For every nonempty B \subseteq A we have F(B) \in B
```

**Proposition 3.1.6** (SC). For any relation R between A and B, there exists a function  $H: A \to B$  such that  $H \subseteq R$  (i.e.  $\forall x \in A.xRH(x)$ ).

```
PROOF: \langle 1 \rangle 1. Let: R be a relation between A and B. \langle 1 \rangle 2. Pick a choice function G for B. \langle 1 \rangle 3. Define H: A \to B by H(x) = G(\{y \mid xRy\}) \langle 1 \rangle 4. H \subseteq R
```

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## 3.1.1 Injective Functions

**Definition 3.1.7** (Injective). A function  $\mathbf{F} : \mathbf{A} \to \mathbf{B}$  is one-to-one, injective or an injection,  $\mathbf{F} : \mathbf{A} \rightarrowtail \mathbf{B}$ , iff, for all  $x, y \in \mathbf{A}$ , if  $\mathbf{F}(x) = \mathbf{F}(y)$ , then x = y.

**Proposition 3.1.8** (S). For any class A, the following is a theorem:  $id_A : A \to A$  is injective.

PROOF: If  $id_{\mathbf{A}}(x) = id_{\mathbf{A}}(y)$  then immediately x = y.  $\square$ 

**Proposition Schema 3.1.9** (S). For any classes **A**, **B**, **C**, **F**, **G**, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \rightarrowtail \mathbf{B}$  and  $\mathbf{G}: \mathbf{B} \rightarrowtail \mathbf{C}$ . Then  $\mathbf{G} \circ \mathbf{F}: \mathbf{A} \rightarrowtail \mathbf{C}$ .

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $x, y \in \mathbf{A}$
- $\langle 1 \rangle 2$ . Assume:  $(\mathbf{G} \circ \mathbf{F})(x) = (\mathbf{G} \circ \mathbf{F})(y)$
- $\langle 1 \rangle 3. \ \mathbf{G}(\mathbf{F}(x)) = \mathbf{G}(\mathbf{F}(y))$
- $\langle 1 \rangle 4$ .  $\mathbf{F}(x) = \mathbf{F}(y)$

PROOF: Since G is injective.

 $\langle 1 \rangle 5. \ x = y$ 

PROOF: Since  $\mathbf{F}$  is injective.

П

**Proposition 3.1.10** (S). Let  $F: A \to B$  where A is nonempty. There exists  $G: B \to A$  (a left inverse) such that  $G \circ F = \mathrm{id}_A$  if and only if F is one-to-one.

## Proof:

- $\langle 1 \rangle 1$ . If there exists  $G: B \to A$  such that  $G \circ F = \mathrm{id}_A$  then F is one-to-one.
  - $\langle 2 \rangle 1$ . Assume:  $G: B \to A$  and  $G \circ F = I_A$
  - $\langle 2 \rangle 2$ . Let:  $x, y \in A$
  - $\langle 2 \rangle 3$ . Assume: F(x) = F(y)
  - $\langle 2 \rangle 4. \ x = y$

PROOF: x = G(F(x)) = G(F(y)) = y

- $\langle 1 \rangle 2$ . If F is one-to-one then there exists  $G: B \to A$  such that  $G \circ F = I_A$ .
  - $\langle 2 \rangle 1$ . Assume: F is one-to-one.
  - $\langle 2 \rangle 2$ . Pick  $a \in A$
  - $\langle 2 \rangle 3$ . Let:  $G: B \to A$  be the function defined by: G(b) is the (unique)  $x \in A$  such that F(x) = b if there exists such an x, G(b) = a otherwise.
  - $\langle 2 \rangle 4$ . For all  $x \in A$  we have G(F(x)) = x.

3.1.2 Surjective Functions

**Definition 3.1.11** (Surjective). Let  $F:A\to B$ . We say that F is *surjective*, or maps A onto B, and write  $F:A\twoheadrightarrow B$ , iff for all  $y\in B$  there exists  $x\in A$  such that F(x)=y.

**Proposition Schema 3.1.12** (S). For any class A, the following is a theorem:  $id_A : A \to A$  is surjective.

PROOF: For any  $y \in \mathbf{A}$  we have  $\mathrm{id}_{\mathbf{A}}(y) = y$ .  $\square$ 

**Proposition Schema 3.1.13** (S). For any classes A, B, C, F and G, the following is a theorem:

If  $\mathbf{F} : \mathbf{A} \twoheadrightarrow \mathbf{B}$  and  $\mathbf{G} : \mathbf{B} \twoheadrightarrow \mathbf{C}$ , then  $\mathbf{G} \circ \mathbf{F} : \mathbf{A} \twoheadrightarrow \mathbf{C}$ .

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $c \in \mathbf{C}$
- $\langle 1 \rangle 2$ . Pick  $b \in \mathbf{B}$  such that  $\mathbf{G}(b) = c$ .
- $\langle 1 \rangle 3$ . Pick  $a \in \mathbf{A}$  such that  $\mathbf{F}(a) = b$ .
- $\langle 1 \rangle 4. \ (\mathbf{G} \circ \mathbf{F})(a) = c$

**Proposition 3.1.14** (SC). Let  $F: A \to B$ . There exists  $H: B \to A$  (a right inverse) such that  $F \circ H = \mathrm{id}_B$  if and only if F maps A onto B.

### Proof:

- $\langle 1 \rangle 1$ . If F has a right inverse then F is surjective.
  - $\langle 2 \rangle 1$ . Assume: F has a right inverse  $H: B \to A$ .
  - $\langle 2 \rangle 2$ . Let:  $y \in B$
  - $\langle 2 \rangle 3$ . F(H(y)) = y
  - $\langle 2 \rangle 4$ . There exists  $x \in A$  such that F(x) = y
- $\langle 1 \rangle 2$ . If F is surjective then F has a right inverse.
  - $\langle 2 \rangle 1$ . Assume: F is surjective.
  - $\langle 2 \rangle 2$ . PICK a function  $H: B \to A$  such that  $H \subseteq F^{-1}$  PROOF: Proposition 3.1.6.
  - $\langle 2 \rangle 3. \ F \circ H = \mathrm{id}_B$ 
    - $\langle 3 \rangle 1$ . Let:  $y \in B$
    - $\langle 3 \rangle 2. \ (y, H(y)) \in F^{-1}$
    - $\langle 3 \rangle 3. \ F(H(y)) = y$

### 3.1.3 Bijections

**Definition 3.1.15** (Bijection). Let  $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ . Then  $\mathbf{F}$  is *bijective* or a *bijection*,  $\mathbf{F} : \mathbf{A} \approx \mathbf{B}$ , iff it is injective and surjective.

**Proposition Schema 3.1.16** (S). For any class A, the following is a theorem: The identity function  $\mathrm{id}_A: A \approx A$  is a bijection.

Proof: Proposition 3.1.8 and 3.1.12.  $\square$ 

**Proposition Schema 3.1.17** (S). For any classes A, B and F, the following is a theorem:

If  $\mathbf{F} : \mathbf{A} \approx \mathbf{B}$  then  $\mathbf{F}^{-1} : \mathbf{B} \approx \mathbf{A}$ .

#### Proof:

- $\langle 1 \rangle 1. \ \mathbf{F}^{-1} : \mathbf{B} \to \mathbf{A}$ 
  - $\langle 2 \rangle 1$ . Let:  $b \in \mathbf{B}$

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 $\langle 2 \rangle 2$ . PICK  $a \in \mathbf{A}$  such that  $\mathbf{F}(a) = b$ .

Proof: Since  $\mathbf{F}$  is surjective.

 $\langle 2 \rangle 3. \ (b,a) \in \mathbf{F}^{-1}$ 

 $\langle 2 \rangle 4$ . If  $(b, a') \in \mathbf{F}^{-1}$  then a' = a.

 $\langle 3 \rangle 1$ . Let:  $a' \in \mathbf{A}$  such that  $(b, a') \in \mathbf{F}^{-1}$ 

 $\langle 3 \rangle 2$ .  $\mathbf{F}(a') = \mathbf{F}(a)$ 

 $\langle 3 \rangle 3. \ a' = a$ 

Proof: Since  $\mathbf{F}$  is injective.

 $\langle 1 \rangle 2$ .  $\mathbf{F}^{-1}$  is injective.

 $\langle 2 \rangle 1$ . Let:  $x, y \in \mathbf{B}$ 

 $\langle 2 \rangle 2$ . Assume:  $\mathbf{F}^{-1}(x) = \mathbf{F}^{-1}(y)$ 

 $\langle 2 \rangle 3. \ x = y$ 

PROOF:  $x = \mathbf{F}(\mathbf{F}^{-1}(x)) = \mathbf{F}(\mathbf{F}^{-1}(y)) = y$ .

 $\langle 1 \rangle 3$ .  $\mathbf{F}^{-1}$  is surjective.

PROOF: For all  $a \in \mathbf{A}$  we have  $\mathbf{F}^{-1}(\mathbf{F}(a)) = a$ .

П

**Proposition Schema 3.1.18** (S). For any classes **A**, **B**, **C**, **F** and **G**, the following is a theorem:

If  $\mathbf{F} : \mathbf{A} \approx \mathbf{B}$  and  $\mathbf{G} : \mathbf{B} \approx \mathbf{C}$  then  $\mathbf{G} \circ \mathbf{F} : \mathbf{A} \approx \mathbf{C}$ .

Proof: Propositions 3.1.9 and 3.1.13.  $\square$ 

#### 3.1.4 Restrictions

**Definition 3.1.19** (Restriction). Let  $\mathbf{F} : \mathbf{A} \to \mathbf{B}$ . Let  $\mathbf{C} \subseteq \mathbf{A}$ . The *restriction* of  $\mathbf{F}$  to  $\mathbf{C}$ , denoted  $\mathbf{F} \upharpoonright \mathbf{C}$ , is the function

$$\mathbf{F} \upharpoonright \mathbf{C} : \mathbf{C} \to \mathbf{B}$$
 
$$(\mathbf{F} \upharpoonright \mathbf{C})(x) = \mathbf{F}(x) \qquad (x \in \mathbf{C})$$

## **3.1.5** Images

**Definition 3.1.20** (Image). Let  $F:A\to B$  and  $C\subseteq A$ . The *image* of C under F is the class

$$\mathbf{F}(\mathbf{C}) := \{ \mathbf{F}(x) \mid x \in \mathbf{C} \} .$$

**Proposition Schema 3.1.21** (S). For any classes **F**, **A** and **B**, the following is a theorem.

If  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$ , then for any subset  $S \subseteq \mathbf{A}$ , the class  $\mathbf{F}(S)$  is a set.

PROOF: By an Axiom of Replacement.

**Proposition Schema 3.1.22** (S). For any classes A, B, C and F, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C} \subseteq \mathcal{P}\mathbf{A}$ . Then

$$\mathbf{F}\left(\bigcup\mathbf{C}\right) = \{y \mid \exists X \in \mathbf{C}.y \in \mathbf{F}(X)\}$$

Proof:

$$y \in \mathbf{F}\left(\bigcup \mathbf{C}\right) \Leftrightarrow \exists x \in \bigcup \mathbf{C}. y = \mathbf{F}(x)$$
  
 $\Leftrightarrow \exists x. \exists X. X \in \mathbf{C} \land x \in X \land y = \mathbf{F}(x)$   
 $\Leftrightarrow \exists X \in \mathbf{C}. y \in \mathbf{F}(X)$ 

**Proposition Schema 3.1.23** (S). For any classes A, B, C, D and F, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C}, \mathbf{D} \subseteq \mathbf{A}$ . Then

$$\mathbf{F}(\mathbf{C} \cup \mathbf{D}) = \mathbf{F}(\mathbf{C}) \cup \mathbf{F}(\mathbf{D})$$
.

Proof:

$$y \in \mathbf{F}(\mathbf{C} \cup \mathbf{D}) \Leftrightarrow \exists x \in \mathbf{C} \cup \mathbf{D}. y = \mathbf{F}(x)$$
  
 $\Leftrightarrow \exists x \in \mathbf{C}. y = \mathbf{F}(x) \lor \exists x \in \mathbf{D}. y = \mathbf{F}(x)$   
 $\Leftrightarrow y \in \mathbf{F}(\mathbf{C}) \cup \mathbf{F}(\mathbf{D})$ 

**Proposition 3.1.24** (S). For any classes **F**, **A**, **B**, **C** and **D**, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C}, \mathbf{D} \subseteq \mathbf{A}$ . Then

$$F(A \cap B) \subseteq F(A) \cap F(B)$$
.

Equality holds if  $\mathbf{F}$  is injective.

Proof:

```
\langle 1 \rangle 1. \mathbf{F}(\mathbf{A} \cap \mathbf{B}) \subseteq \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B})
     \langle 2 \rangle 1. Let: y \in \mathbf{F}(\mathbf{A} \cap \mathbf{B})
     \langle 2 \rangle 2. Pick x \in \mathbf{A} \cap \mathbf{B} such that y = \mathbf{F}(x)
     \langle 2 \rangle 3. \ y \in \mathbf{F}(\mathbf{A})
          PROOF: Since x \in \mathbf{A}.
     \langle 2 \rangle 4. \ y \in \mathbf{F}(\mathbf{B})
          PROOF: Since x \in \mathbf{B}.
\langle 1 \rangle 2. If F is injective then \mathbf{F}(\mathbf{A} \cap \mathbf{B}) = \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B}).
     \langle 2 \rangle 1. Assume: F is injective.
     \langle 2 \rangle 2. Let: y \in \mathbf{F}(\mathbf{A}) \cap \mathbf{F}(\mathbf{B})
     \langle 2 \rangle 3. Pick x \in \mathbf{A} such that y = \mathbf{F}(x)
     \langle 2 \rangle 4. Pick x' \in \mathbf{B} such that y = \mathbf{F}(x')
     \langle 2 \rangle 5. \ x = x'
          Proof: \langle 2 \rangle 1
     \langle 2 \rangle 6. \ x \in \mathbf{A} \cap \mathbf{B}
     \langle 2 \rangle 7. \ y \in \mathbf{F}(\mathbf{A} \cap \mathbf{B})
```

**Proposition Schema 3.1.25** (S). For any classes **F**, **A**, **B**, and **C**, the following is a theorem:

Let  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C} \subseteq \mathcal{P}\mathbf{A}$ . Then

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$$\mathbf{F}\left(\bigcap \mathbf{C}\right) \subseteq \bigcap \{\mathbf{F}(X) \mid X \in \mathbf{A}\}$$
.

Equality holds if **F** is injective and **A** is nonempty.

```
Proof:
```

```
\langle 1 \rangle 1. \ \mathbf{F} (\bigcap \mathbf{A}) \subseteq \bigcap \{ \mathbf{F}(X) \mid X \in \mathbf{A} \}
     \langle 2 \rangle 1. Let: y \in \mathbf{F}(\bigcap \mathbf{A})
     \langle 2 \rangle 2. PICK x \in \bigcap \mathbf{A} such that y = \mathbf{F}(x)
     \langle 2 \rangle 3. Let: X \in \mathbf{A}
     \langle 2 \rangle 4. \ x \in X
     \langle 2 \rangle 5. \ y \in \mathbf{F}(X)
\langle 1 \rangle 2. If F is injective then \mathbf{F}(\bigcap \mathbf{A}) = \bigcap \{ \mathbf{F}(X) \mid X \in \mathbf{A} \}
     \langle 2 \rangle 1. Assume: F is injective.
     \langle 2 \rangle 2. Assume: A is nonempty.
     \langle 2 \rangle 3. Let: y \in \bigcap \{ \mathbf{F}(X) \mid X \in \mathbf{A} \}
     \langle 2 \rangle 4. Pick X_0 \in \mathbf{A}
     \langle 2 \rangle5. Pick x \in X_0 such that (x, y) \in \mathbf{F}
     \langle 2 \rangle 6. \ x \in \bigcap \mathbf{A}
          \langle 3 \rangle 1. Let: X \in \mathbf{A}
          \langle 3 \rangle 2. PICK x' \in X such that (x', y) \in \mathbf{F}.
          \langle 3 \rangle 3. \ x = x'
               Proof: \langle 2 \rangle 1
           \langle 3 \rangle 4. \ x \in X
     \langle 2 \rangle 7. \ y \in \mathbf{F}(\bigcap \mathbf{A})
```

**Proposition 3.1.26** (S). For any classes A, B, C, D and F, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C}, \mathbf{D} \subseteq \mathbf{A}$ . Then

$$\mathbf{F}(\mathbf{C}) - \mathbf{F}(\mathbf{D}) \subseteq \mathbf{F}(\mathbf{C} - \mathbf{D}) \ .$$

Equality holds if  $\mathbf{F}$  is injective.

Proof:

$$\begin{array}{l} \langle 1 \rangle 1. \ \mathbf{F(C)} - \mathbf{F(D)} \subseteq \mathbf{F(A-B)} \\ \langle 2 \rangle 1. \ \mathrm{LET:} \ y \in \mathbf{F(A)} - \mathbf{F(B)} \\ \langle 2 \rangle 2. \ \mathrm{PICK} \ x \in \mathbf{A} \ \mathrm{such \ that} \ y = \mathbf{F}(x) \\ \langle 2 \rangle 3. \ x \notin \mathbf{B} \\ \langle 2 \rangle 4. \ x \in \mathbf{A} - \mathbf{B} \\ \langle 2 \rangle 5. \ y \in \mathbf{F(A-B)} \\ \langle 1 \rangle 2. \ \mathrm{If} \ \mathbf{F} \ \mathrm{is} \ \mathrm{injective} \ \mathrm{then} \ \mathbf{F(A)} - \mathbf{F(B)} = \mathbf{F(A-B)} \\ \langle 2 \rangle 1. \ \mathrm{Assume:} \ \mathbf{F} \ \mathrm{is} \ \mathrm{injective}. \\ \langle 2 \rangle 2. \ \mathrm{Let:} \ y \in \mathbf{F(A-B)} \\ \langle 2 \rangle 3. \ \mathrm{PICK} \ x \in \mathbf{A} - \mathbf{B} \ \mathrm{such \ that} \ y = \mathbf{F}(x) \\ \langle 2 \rangle 4. \ y \in \mathbf{F(A)} \\ \langle 2 \rangle 5. \ y \notin \mathbf{F(B)} \end{array}$$

- $\langle 3 \rangle 1$ . Assume: for a contradiction  $y \in \mathbf{F}(\mathbf{B})$
- $\langle 3 \rangle 2$ . Pick  $x' \in \mathbf{B}$  such that  $y = \mathbf{F}(x')$
- $\langle 3 \rangle 3. \ x = x'$

Proof:  $\langle 2 \rangle 1$ 

- $\langle 3 \rangle 4. \ x \in \mathbf{B}$
- $\langle 3 \rangle 5$ . Q.E.D.

Proof: This contradicts  $\langle 2 \rangle 3$ .

## 3.1.6 Inverse Images

**Definition 3.1.27** (Inverse Image). Let  $F:A\to B$  and  $C\subseteq B$ . Then the *inverse image* of C under F is

$$\mathbf{F}^{-1}(\mathbf{C}) = \{x \in \mathbf{A} \mid \mathbf{F}(x) \in \mathbf{C}\}\$$
.

**Proposition Schema 3.1.28** (S). For any classes A, B, C and F, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C} \subseteq \mathcal{P}\mathbf{B}$ . Then

$$\mathbf{F}^{-1}\left(\bigcap \mathbf{C}\right) = \bigcap \{\mathbf{F}^{-1}(X) \mid X \in \mathbf{C}\}\ .$$

Proof:

$$x \in \mathbf{F}^{-1}\left(\bigcap \mathbf{C}\right) \Leftrightarrow \mathbf{F}(x) \in \bigcap \mathbf{C}$$
  
 $\Leftrightarrow \forall X \in \mathbf{C}.\mathbf{F}(x) \in X$   
 $\Leftrightarrow \forall X \in \mathbf{C}.x \in \mathbf{F}^{-1}(X)$ 

**Proposition Schema 3.1.29** (S). For any classes A, B, C, D and F, the following is a theorem:

Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{C}, \mathbf{D} \subseteq \mathbf{B}$ . Then

$$\mathbf{F}^{-1}(\mathbf{C}-\mathbf{D}) = \mathbf{F}^{-1}(\mathbf{C}) - \mathbf{F}^{-1}(\mathbf{D}) \ .$$

Proof:

$$x \in \mathbf{F}^{-1}(\mathbf{C} - \mathbf{D}) \Leftrightarrow \mathbf{F}(x) \in \mathbf{C} - \mathbf{D}$$
  
 $\Leftrightarrow \mathbf{F}(x) \in \mathbf{C} \wedge \mathbf{F}(x) \notin \mathbf{D}$   
 $\Leftrightarrow x \in \mathbf{F}^{-1}(\mathbf{C}) \wedge x \in \mathbf{F}^{-1}(\mathbf{D})$   
 $\Leftrightarrow x \in \mathbf{F}^{-1}(\mathbf{C}) - \mathbf{F}^{-1}(\mathbf{D})$ 

### 3.1.7 Function Sets

**Proposition 3.1.30** (SF). For any classes **B** and **F**, the following is a theorem: Let A be a set. If  $F : A \to B$  then F is a set.

PROOF: By an Axiom of Replacement, we have  $R = \{ \mathbf{F}(x) \mid x \in A \}$  is a set. Hence  $\mathbf{F}$  is a set since  $\mathbf{F} \subseteq A \times R$ .  $\square$ 

**Definition 3.1.31** (Dependent Product Class). Let I be a set and let  $\mathbf{H}(i)$  be a class for all  $i \in I$ . We write  $\prod_{i \in I} \mathbf{H}(i)$  for the class of all functions  $f: I \to \bigcup_{i \in I} \mathbf{H}(i)$  such that  $\forall i \in I. f(i) \in \mathbf{H}(i)$ . We write  $\mathbf{B}^I$  for  $\prod_{i \in I} \mathbf{B}$  where  $\mathbf{B}$  does not depend on I.

**Proposition Schema 3.1.32** (SF). Let I be a set. Let H(i) be a set for every  $i \in I$ . Then  $\prod_{i \in I} H(i)$  is a set.

```
PROOF: \langle 1 \rangle 1. \{ \mathbf{H}(i) \mid i \in I \} is a set.
PROOF: By an Axiom of Replacement. \langle 1 \rangle 2. \bigcup_{i \in I} \mathbf{H}(i) is a set. \langle 1 \rangle 3. \prod_{i \in I} \mathbf{H}(i) is a set.
PROOF: It is a subset of \mathcal{P}\left(I \times \bigcup_{i \in I} \mathbf{H}(i)\right).
```

**Proposition 3.1.33** (SFC). Let I be a set. Let H(i) be a set for all  $i \in I$ . If  $\forall i \in I. H(i) \neq \emptyset$  then  $\prod_{i \in I} H(i) \neq \emptyset$ .

```
PROOF:  \begin{array}{l} \langle 1 \rangle 1. \text{ Assume: } \forall i \in I.H(i) \neq \emptyset \\ \langle 1 \rangle 2. \text{ Let: } R = \{(i,x) \mid i \in I, x \in H(i)\} \\ \langle 1 \rangle 3. \text{ Pick a function } f:I \rightarrow \bigcup_{i \in I} H(i) \text{ such that } f \subseteq R \\ \text{PROOF: Proposition } 3.1.6. \\ \langle 1 \rangle 4. \ f \in \prod_{i \in I} H(i) \\ \sqcap \end{array}
```

# 3.2 Equinumerosity

**Definition 3.2.1** (Equinumerous). Sets A and B are equinumerous,  $A \approx B$ , iff there exists a bijection between A and B.

## 3.3 Domination

**Definition 3.3.1** (Dominate). A set A is dominated by a set B,  $A \leq B$ , iff there exists an injection  $A \rightarrow B$ .

**Proposition 3.3.2** (SC). Given sets A and B, if  $A \neq \emptyset$  or  $B = \emptyset$ , then we have  $A \leq B$  iff there exists a surjective function  $B \rightarrow A$ .

#### Proof:

- $\langle 1 \rangle 1$ . If  $A \leq B$  and  $A \neq \emptyset$  then there exists a surjective function  $B \to A$ .
  - $\langle 2 \rangle 1$ . Assume:  $f: A \to B$  be injective.
  - $\langle 2 \rangle 2$ . Pick  $a \in A$
  - $\langle 2 \rangle$ 3. Let:  $g: B \to A$  be the function defined by  $g(b) = f^{-1}(b)$  if  $b \in \operatorname{ran} f$ , and g(b) = a otherwise.

```
\langle 2 \rangle 4. g is surjective.
```

 $\langle 1 \rangle 2$ . If there exists a surjective function  $B \to A$  then  $A \leq B$ .

- $\langle 2 \rangle 1$ . Assume: there exists a surjective function  $g: B \to A$
- $\langle 2 \rangle 2$ .  $\forall a \in A. \exists b \in B. g(b) = a$  $\langle 2 \rangle 3$ . Choose a function  $f: A \to B$  such that  $\forall a \in A. g(f(a)) = a$
- $\langle 2 \rangle 4$ . f is injective.

## Chapter 4

# Category Theory

## 4.1 Categories

**Definition 4.1.1** (Category). A category C consists of:

- a class of *objects*;
- for any objects X, Y, a set  $\mathbf{C}(X, Y)$  whose elements are called *morphisms*. We write  $f: X \to Y$  for  $f \in \mathbf{C}(X, Y)$ .
- for any morphisms  $f: X \to Y$  and  $g: Y \to Z$ , a morphism  $gf: X \to Z$

such that:

- For all  $f: X \to Y$ ,  $g: Y \to Z$ ,  $h: Z \to W$ , we have h(gf) = (hg)f.
- For every object X, there exists a morphism  $id_X: X \to X$  such that:
  - for any object Y and morphism  $f: X \to Y$  we have  $fid_X = f$ .
  - for any object Y and morphism  $f: Y \to X$  we have  $id_X f = f$ .

**Definition 4.1.2** (Category of Sets). Let Set be the category of sets and functions.

**Definition 4.1.3** (Category of Pointed Sets). The *category of pointed sets*  $Set_*$  has:

- objects all pairs (A, a) where A is a set and  $a \in A$ ;
- morphisms  $f:(A,a)\to (B,b)$  all functions  $f:A\to B$  such that f(a)=b.

**Definition 4.1.4** (Opposite Category). For any category  $\mathbb{C}$ , the *opposite category*  $\mathbb{C}^{op}$  is the category with the same objects as  $\mathbb{C}$  and  $\mathbb{C}^{op}(X,Y) = \mathbb{C}(Y,X)$ .

## 4.2 Invertible Morphisms

**Definition 4.2.1** (Left Inverse, Right Inverse). In any category, let  $f: A \to B$  and  $g: B \to A$ . Then g is a *left inverse* of f, and f is a *right inverse* of f, iff  $gf = \mathrm{id}_A$ .

**Proposition 4.2.2** (FOL). Let  $f: A \to B$  and  $g, h: B \to A$ . If g is a left inverse to f and h is a right inverse to f then g = h.

PROOF: Since  $g = gid_B = gfh = id_A h = h$ .  $\square$ 

**Definition 4.2.3** (Isomorphism). A morphism  $f: A \to B$  is an *isomorphism*,  $f: A \cong B$ , iff it has both a left and a right inverse. In this case, its unique inverse is denoted  $f^{-1}$ .

Two objects A and B are  $isomorphic,\,A\cong B,$  iff there exists an isomorphism between them.

**Proposition 4.2.4** (S). A function is an isomorphism in Set iff it is a bijection.

**Proposition 4.2.5** (S). Let  $f: A \to B$  in the category  $\mathbb{C}$ . Then the following are equivalent.

- 1. f is an isomorphism.
- 2. For every object X, the function  $\mathbf{C}(X,f):\mathbf{C}(X,A)\to\mathbf{C}(X,B)$  is a bijection.
- 3. For every object X, the function  $\mathbf{C}(f,X):\mathbf{C}(B,X)\to\mathbf{C}(A,X)$  is a bijection.

## Chapter 5

# **Equivalence Relations**

**Definition 5.0.1** (Equivalence Relation). An *equivalence relation* on a class **A** is a binary relation on **A** that is reflexive, symmetric and transitive.

**Proposition 5.0.2** (Z). Equinumerosity is an equivalence relation on the class of all sets.

PROOF: Propositions 3.1.16, 3.1.17, 3.1.18.

**Definition 5.0.3** (Respects). Let **R** be an equivalence relation on **A** and **F**:  $\mathbf{A} \to \mathbf{B}$ . Then **F** respects **A** iff, whenever  $(x,y) \in \mathbf{R}$ , then  $\mathbf{F}(x) = \mathbf{F}(y)$ .

**Definition 5.0.4** (Equivalence Class). Let  $\mathbf{R}$  be an equivalence relation on  $\mathbf{A}$  and  $a \in \mathbf{A}$ . The *equivalence class* of a modulo  $\mathbf{R}$  is

$$[a]_{\mathbf{R}} := \{x \mid a\mathbf{R}x\} .$$

**Proposition Schema 5.0.5** (Z). For any classes  ${\bf A}$  and  ${\bf R}$ , the following is a theorem.

Assume **R** be an equivalence relation on **A**. Let  $a, b \in \mathbf{A}$ . Then  $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$  if and only if  $a\mathbf{R}b$ .

### Proof:

- $\langle 1 \rangle 1$ . If  $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$  then  $a\mathbf{R}b$ .
  - $\langle 2 \rangle 1$ . Assume:  $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$
  - $\langle 2 \rangle 2$ .  $b\mathbf{R}b$

PROOF: Reflexivity

- $\langle 2 \rangle 3. \ b \in [b]_{\mathbf{R}}$
- $\langle 2 \rangle 4. \ b \in [a]_{\mathbf{R}}$
- $\langle 2 \rangle 5$ .  $a\mathbf{R}b$
- $\langle 1 \rangle 2$ . If  $a\mathbf{R}b$  then  $[a]_{\mathbf{R}} = [b]_{\mathbf{R}}$ .
  - $\langle 2 \rangle 1$ . For all  $x, y \in \mathbf{A}$ , if  $x \mathbf{R} y$  then  $[y]_{\mathbf{R}} \subseteq [x]_{\mathbf{R}}$ 
    - $\langle 3 \rangle 1$ . Let:  $x, y \in \mathbf{A}$
    - $\langle 3 \rangle 2$ . Assume:  $x \mathbf{R} y$

```
\langle 3 \rangle 3. \text{ Let: } t \in [y]_{\mathbf{R}}
\langle 3 \rangle 4. y\mathbf{R}t
\langle 3 \rangle 5. x\mathbf{R}t
\text{Proof: Transitivity, } \langle 3 \rangle 2, \langle 3 \rangle 4.
\langle 3 \rangle 6. t \in [x]_{\mathbf{R}}
\langle 2 \rangle 2. \text{ Assume: } a\mathbf{R}b
\langle 2 \rangle 3. [b]_{\mathbf{R}} \subseteq [a]_{\mathbf{R}}
\text{Proof: } \langle 2 \rangle 1, \langle 2 \rangle 2.
\langle 2 \rangle 4. b\mathbf{R}a
\text{Proof: Symmetry, } \langle 2 \rangle 2.
\langle 2 \rangle 5. [a]_{\mathbf{R}} \subseteq [b]_{\mathbf{R}}
\text{Proof: } \langle 2 \rangle 1, \langle 2 \rangle 4.
\langle 2 \rangle 6. [a]_{\mathbf{R}} = [b]_{\mathbf{R}}
\text{Proof: } \langle 2 \rangle 3, \langle 2 \rangle 5.
```

**Definition 5.0.6** (Partition). A partition  $\Pi$  of a set A is a set of nonempty subsets of A that is disjoint and exhaustive, i.e.

- 1. no two different sets in  $\Pi$  have any common elements, and
- 2. each element of A is in some set in  $\Pi$ .

**Definition 5.0.7.** Let R be an equivalence relation on a set A. The *quotient* set A/R is the set of all equivalence classes.

**Theorem 5.0.8** (Z). Let A be a set and  $\mathbf{B}$  a class. Let R be an equivalence relation on A and  $F:A\to \mathbf{B}$ . Then F respects R if and only if there exists  $\hat{F}:A/R\to \mathbf{B}$  such that

$$\forall a \in A. \hat{F}([a]_R) = F(a) .$$

In this case,  $\hat{F}$  is unique.

#### Proof:

- $\langle 1 \rangle 1$ . If F respects R then there exists  $\hat{F}: A/R \to \mathbf{B}$  such that  $\forall a \in A.\hat{F}([a]_R) = F(a)$ .
  - $\langle 2 \rangle 1$ . Assume: F respects R.
  - $\langle 2 \rangle 2$ . Let:  $\hat{F} = \{ ([a]_R, F(a)) \mid a \in A \}$
  - $\langle 2 \rangle 3$ .  $\hat{F}$  is a function.
    - $\langle 3 \rangle 1$ . Assume:  $a, a' \in A$  and  $[a]_R = [a']_R$ Prove: F(a) = F(a')
    - $\langle 3 \rangle 2. \ (a, a') \in R$

Proof: Proposition 5.0.5.

- $\langle 3 \rangle 3. \ F(a) = F(a')$ 
  - Proof:  $\langle 2 \rangle 1$
- $\langle 2 \rangle 4$ . dom  $\hat{F} = A/R$
- $\langle 2 \rangle 5$ . ran  $\hat{F} \subseteq \mathbf{B}$

```
\langle 2 \rangle 6. \ \forall a \in A. \hat{F}([a]_R) = F(a)
\langle 1 \rangle 2. If there exists \hat{F}: A/R \to \mathbf{B} such that \forall a \in A.\hat{F}([a]_R) = F(a) then F
        respects R.
   \langle 2 \rangle 1. Assume: \hat{F}: A/R \to \mathbf{B} and \forall a \in A.\hat{F}([a]_R) = F(a)
   \langle 2 \rangle 2. Let: a, a' \in A
   \langle 2 \rangle 3. Assume: (a, a') \in R
   \langle 2 \rangle 4. [a]_R = [a']_R
      Proof: Proposition 5.0.5.
   \langle 2 \rangle 5. F(a) = F(a')
      Proof: \langle 2 \rangle 1
\langle 1 \rangle 3. If G, H : A/R \to \mathbf{B} and \forall a \in A.G([a]_R) = H([a]_R) then G = H.
Proposition 5.0.9 (Z). Let R be an equivalence relation on a set A. Then
A/R is a partition of A.
Proof:
\langle 1 \rangle 1. Every member of A/R is nonempty.
   PROOF: Since a \in [a]_R by reflexivity.
\langle 1 \rangle 2. No two different sets in A/R have any common elements.
   \langle 2 \rangle 1. Let: [a]_R, [b]_R \in A/R
   \langle 2 \rangle 2. Let: c \in [a]_R \cap [b]_R
            PROVE: [a]_R = [b]_R
   \langle 2 \rangle 3. aRc
      Proof: \langle 2 \rangle 2
   \langle 2 \rangle 4. \ bRc
      Proof: \langle 2 \rangle 2
   \langle 2 \rangle 5. cRb
      Proof: Symmetry, \langle 2 \rangle 4
   \langle 2 \rangle 6. aRb
      Proof: Transitivity, \langle 2 \rangle 3, \langle 2 \rangle 5
   \langle 2 \rangle 7. [a]_R = [b]_R
      Proof: Proposition 5.0.5, \langle 2 \rangle 6
\langle 1 \rangle 3. Each element of A is in some set in A/R.
   PROOF: Since a \in [a]_R by reflexivity.
```

**Proposition 5.0.10** (Z). For any partition P of a set A, there exists a unique equivalence relation R on A such that A/R = P, namely xRy iff  $\exists X \in P(x \in X \land y \in X)$ .

Proof: Easy.

**Definition 5.0.11** (Natural Map). Let A be a set and R an equivalence relation on A. The natural map  $A \to A/R$  is the function that maps  $a \in A$  to  $[a]_R$ .

## Chapter 6

# Ordering Relations

## 6.1 Partial Orders

**Definition 6.1.1** (Partial Ordering). Let **A** be a class. A *partial ordering* on **A** is a relation **R** on **A** that is reflexive, antisymmetric and transitive.

We often write  $\leq$  for a partial ordering, and then write x < y for  $x \leq y \land x \neq y$ .

**Proposition Schema 6.1.2** (Z). For any classes A and R, the following is a theorem:

If **R** is a partial order on **A** then so is  $\mathbf{R}^{-1}$ .

```
Proof:
```

```
\langle 1 \rangle 1. \mathbf{R}^{-1} is reflexive.

PROOF: Proposition 2.2.12.
\langle 1 \rangle 2. \mathbf{R}^{-1} is antisymmetric.

PROOF: Proposition 2.2.16.
\langle 1 \rangle 3. \mathbf{R}^{-1} is transitive.
\langle 2 \rangle 1. ASSUME: x\mathbf{R}^{-1}y and y\mathbf{R}^{-1}z
\langle 2 \rangle 2. y\mathbf{R}x and z\mathbf{R}y
\langle 2 \rangle 3. z\mathbf{R}x

PROOF: Since \mathbf{R} is transitive.
\langle 2 \rangle 4. x\mathbf{R}^{-1}z
```

**Proposition Schema 6.1.3** (Z). For any classes A, B, F and R, the following is a theorem:

Assume **R** is a partial order on **B** and  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  is injective. Define **S** on **A** by  $x\mathbf{S}y$  iff  $\mathbf{F}(x)\mathbf{RF}(y)$ . Then **S** is a partial order on **A**.

#### Proof:

 $\langle 1 \rangle 1$ . **S** is reflexive.

PROOF: For any  $x \in \mathbf{A}$  we have  $\mathbf{F}(x)\mathbf{RF}(x)$ .

```
\langle 1 \rangle2. S is antisymmetric.

\langle 2 \rangle1. Let: x, y \in \mathbf{A}

\langle 2 \rangle2. Assume: x\mathbf{S}y and y\mathbf{S}x

\langle 2 \rangle3. \mathbf{F}(x)\mathbf{R}\mathbf{F}(y) and \mathbf{F}(y)\mathbf{R}\mathbf{F}(x)

\langle 2 \rangle4. \mathbf{F}(x) = \mathbf{F}(y)

PROOF: R is antisymmetric.

\langle 2 \rangle5. x = y

\langle 1 \rangle3. S is transitive.
```

Corollary Schema 6.1.3.1 (Z). For any classes A, B and R, the following is a theorem:

Assume **R** be a partial order on **A** and **B**  $\subseteq$  **A**. Then **R**  $\cap$  **B**<sup>2</sup> is a partial order on **B**.

**Definition 6.1.4** (Partially Ordered Set). A partially ordered set or poset is a pair  $(A, \leq)$  where A is a set and  $\leq$  is a partial ordering on A. We often write just A for  $(A, \leq)$ .

If  $(A, \leq)$  is a poset and  $B \subseteq A$  we write just B for the poset  $(B, \leq \cap B^2)$ .

**Definition 6.1.5** (Strictly Monotone). Let  $(A, <_A)$  and  $(B, <_B)$  be posets. A function  $f: A \to B$  is *strictly monotone* iff, whenever  $x <_A y$ , then  $f(x) <_B f(y)$ .

**Definition 6.1.6** (Least). Let  $\leq$  be a partial order on  $\mathbf{A}$ . An element  $m \in \mathbf{A}$  is *least* iff for all  $x \in \mathbf{A}$  we have  $m \leq x$ .

**Proposition 6.1.7** (Z). A partial order has at most one least element.

PROOF: If m and m' are least then  $m \leq m'$  and  $m' \leq m$ , so m = m'.  $\square$ 

**Definition 6.1.8** (Greatst). Let  $\leq$  be a partial order on **A**. An element  $m \in \mathbf{A}$  is *greatest* iff for all  $x \in A$  we have  $x \leq m$ .

**Proposition 6.1.9** (Z). A poset has at most one greatest element.

PROOF: If m and m' are greatest then  $m \leq m'$  and  $m' \leq m$ , so m = m'.  $\square$ 

**Definition 6.1.10** (Upper Bound). Let  $\leq$  be a partial ordering on **A** and  $\mathbf{B} \subseteq \mathbf{A}$ . Let  $u \in \mathbf{A}$ . Then u is an *upper bound* for **B** iff  $\forall x \in \mathbf{B}.x \leq u$ .

**Definition 6.1.11** (Lower Bound). Let  $\leq$  be a partial ordering on **A** and  $\mathbf{B} \subseteq \mathbf{A}$ . Let  $l \in \mathbf{A}$ . Then l is a *lower bound* for **B** iff  $\forall x \in \mathbf{B}.l \leq x$ .

**Definition 6.1.12** (Bounded Above). Let  $\leq$  be a partial ordering on **A** and  $\mathbf{B} \subseteq \mathbf{A}$ . Then **B** is *bounded above* iff it has an upper bound.

**Definition 6.1.13** (Bounded Below). Let  $\leq$  be a partial ordering on **A** and  $\mathbf{B} \subseteq \mathbf{A}$ . Then **B** is *bounded below* iff it has a lower bound.

**Definition 6.1.14** (Least Upper Bound). Let  $\leq$  be a partial ordering on  $\mathbf{A}$  and  $\mathbf{B} \subseteq \mathbf{A}$ . Let  $s \in \mathbf{A}$ . Then s is the *least upper bound* or *supremum* of  $\mathbf{B}$  iff s is an upper bound for  $\mathbf{B}$  and, for every upper bound u for  $\mathbf{B}$ , we have  $s \leq u$ .

**Definition 6.1.15** (Greatest Lower Bound). Let  $\leq$  be a partial ordering on  $\mathbf{A}$  and  $\mathbf{B} \subseteq \mathbf{A}$ . Let  $i \in \mathbf{A}$ . Then i is the *greatest lower bound* or *infimum* of  $\mathbf{B}$  iff i is a lower bound for  $\mathbf{B}$  and, for every lower bound l for  $\mathbf{B}$ , we have  $i \leq l$ .

**Definition 6.1.16** (Complete). A partial order is *complete* iff every nonempty subset bounded above has a supremum, and every nonempty subset bounded below has an infimum.

**Definition 6.1.17** (Order Isomorphism). Let A and B be posets. An order isomorphism between A and B,  $f:A\cong B$ , is a bijection  $f:A\approx B$  such that f and  $f^{-1}$  are monotone.

**Theorem 6.1.18** (Knaster Fixed-Point Theorem (Z)). Let A be a complete poset with a greatest and least element. Let  $\phi: A \to A$  be monotone. Then there exists  $a \in A$  such that  $\phi(a) = a$ .

#### Proof:

```
\langle 1 \rangle 1. Let: B = \{ x \in A \mid x \le \phi(x) \}
```

 $\langle 1 \rangle 2$ . Let:  $a = \sup B$ 

PROOF: B is nonempty because the least element of A is in B, and it is bounded above by the greatest element of A.

```
\langle 1 \rangle 3. For all b \in B we have b \leq \phi(a)
```

```
\langle 2 \rangle 1. Let: b \in B
```

$$\langle 2 \rangle 2. \ b \le \phi(b)$$

$$\langle 2 \rangle 3. \ b \leq a$$

$$\langle 2 \rangle 4. \ \phi(b) \leq \phi(a)$$

$$\langle 2 \rangle 5.$$
  $b \leq \phi(a)$ 

$$\langle 1 \rangle 4. \ a \leq \phi(a)$$

$$\langle 1 \rangle 5. \ \phi(a) \le \phi(\phi(a))$$

$$\langle 1 \rangle 6. \ \phi(a) \in B$$

$$\langle 1 \rangle 7. \ \phi(a) \leq a$$

$$\langle 1 \rangle 8. \ \phi(a) = a$$

**Definition 6.1.19** (Dense). Let  $\leq$  be a partial order on **A** and **B**  $\subseteq$  **A**. Then **B** is *dense* iff, for all  $x, y \in$  **A**, if x < y then there exists  $z \in$  **B** such that x < z < y.

**Proposition 6.1.20** (Z). Let A be a complete poset with no least element. Let  $B \subseteq A$  be dense. Let  $\theta : A \to A$  be a monotone map that is the identity on B. Then  $\theta = id_A$ .

```
\langle 1 \rangle 1. Let: a \in A
Prove: \theta(a) = a
```

```
\langle 1 \rangle 2. Let: S(a) = \{ b \in B \mid b < a \}
\langle 1 \rangle 3. S(a) is nonempty and bounded above.
   \langle 2 \rangle 1. S(a) is nonempty.
      \langle 3 \rangle 1. Pick a_1 < a
          Proof: Since a is not least.
       \langle 3 \rangle 2. There exists b \in B such that a_1 < b < a.
   \langle 2 \rangle 2. S(a) is bounded above by a.
\langle 1 \rangle 4. sup S(a) \leq a
\langle 1 \rangle 5. sup S(a) = a
   \langle 2 \rangle 1. Assume: for a contradiction sup S(a) < a
   \langle 2 \rangle 2. Pick b \in B such that \sup S(a) < b < a
   \langle 2 \rangle 3. \ b \in S(a)
   \langle 2 \rangle 4. Q.E.D.
      PROOF: This contradicts the fact that \sup S(a) < b.
\langle 1 \rangle 6. For all b \in S(a) we have b \leq \theta(a)
   \langle 2 \rangle 1. Let: b \in S(a)
   \langle 2 \rangle 2. b < a
   \langle 2 \rangle 3. \ \theta(b) \leq \theta(a)
   \langle 2 \rangle 4. \ b \leq \theta(a)
      Proof: \theta(b) = b
\langle 1 \rangle 7. \ a \leq \theta(a)
  PROOF: Since a = \sup S(a) and \theta(a) is an upper bound for S(a).
\langle 1 \rangle 8. \ a \not< \theta(a)
   \langle 2 \rangle 1. Assume: for a contradiction a < \theta(a).
   \langle 2 \rangle 2. Pick b \in B such that a < b < \theta(a)
   \langle 2 \rangle 3. \theta(a) \leq \theta(b) = b
   \langle 2 \rangle 4. Q.E.D.
      PROOF: This contradicts the fact that b < \theta(a).
\langle 1 \rangle 9. \ \theta(a) = a
```

**Theorem 6.1.21** (Z). Let A and P be complete posets with no least or greatest element. Let B be dense in A and Q be dense in P. Every order isomorphism  $\phi: B \cong Q$  extends uniquely to an order isomorphism  $A \cong P$ .

#### Proof:

 $\langle 1 \rangle 1$ . For  $a \in A$ , let  $S(a) = \{b \in B \mid b < a\}$ .  $\langle 1 \rangle 2$ . Define  $\overline{\phi} : A \to P$  by  $\overline{\phi}(a) = \sup \phi(S(a))$ .  $\langle 2 \rangle 1$ .  $\phi(S(a))$  is nonempty.  $\langle 3 \rangle 1$ . PICK  $a_1 < a$ PROOF: Since a is not least.  $\langle 3 \rangle 2$ . PICK  $b \in B$  such that  $a_1 < b < a$ .  $\langle 3 \rangle 3$ .  $\phi(b) \in \phi(S(a))$   $\langle 2 \rangle 2$ .  $\phi(S(a))$  is bounded above.  $\langle 3 \rangle 1$ . PICK  $a_2 > a$ PROOF: Since a is not greatest.

 $\langle 3 \rangle 2$ . Pick  $b \in B$  such that  $a < b < a_2$ 

```
\langle 3 \rangle 3. \phi(b) is an upper bound for \phi(S(a)).
\langle 1 \rangle 3. \overline{\phi} is monotone.
   PROOF: If a \leq a' then S(a) \subseteq S(a') and so \overline{\phi}(a) \leq \overline{\phi}(a').
\langle 1 \rangle 4. \phi extends \phi.
   \langle 2 \rangle 1. Let: b \in B
             PROVE: \phi(b) = \sup \phi(S(b))
    \langle 2 \rangle 2. \phi(b) is an upper bound for \phi(S(b))
   \langle 2 \rangle 3. Let: u be any upper bound for \phi(S(b))
             Prove: \phi(b) < u
   \langle 2 \rangle 4. Assume: for a contradiction u < \phi(b)
   \langle 2 \rangle5. Pick q \in Q such that u < q < \phi(b)
   \langle 2 \rangle 6. Pick b' \in B such that \phi(b') = q
    \langle 2 \rangle 7. \ b' < b
    \langle 2 \rangle 8. \ b' \in S(b)
    \langle 2 \rangle 9. \ \ q = \phi(b') \leq u
   \langle 2 \rangle 10. Q.E.D.
       Proof: This is a contradiction.
\langle 1 \rangle 5. Let: \psi = \phi^{-1}
(1)6. Let: \overline{\psi}: P \to A be the function \overline{\psi}(p) = \sup\{\psi(q) \mid q \in Q, q < p\}
\langle 1 \rangle 7. \overline{\psi} is monotone and extends \psi
   PROOF: Similar.
\langle 1 \rangle 8. \overline{\psi} \circ \overline{\phi} : A \to A is monotone and the identity on B.
\langle 1 \rangle 9. \ \overline{\psi} \circ \overline{\phi} = \mathrm{id}_A
   Proof: Proposition 6.1.20.
\langle 1 \rangle 10. \ \overline{\phi} \circ \overline{\psi} = \mathrm{id}_B
   Proof: Proposition 6.1.20.
\langle 1 \rangle 11. If \phi^* : A \cong P is any order isomorphism that extends \phi then \phi^* = \overline{\phi}.
   \langle 2 \rangle 1. Let: a \in A
             PROVE: \phi^*(a) = \sup \phi(S(a))
    \langle 2 \rangle 2. \phi^*(a) is an upper bound for \phi(S(a))
   \langle 2 \rangle 3. Let: u be any upper bound for \phi(S(a))
             PROVE: \phi^*(a) \le u
    \langle 2 \rangle 4. Assume: for a contradiction u < \phi^*(a)
    \langle 2 \rangle5. Pick q \in Q such that u < q < \phi^*(a)
    \langle 2 \rangle 6. Pick b \in B such that q = \phi(b)
    \langle 2 \rangle 7. b < a
    \langle 2 \rangle 8. \ b \in S(a)
   \langle 2 \rangle 9. \ q = \phi(b) \leq u
   \langle 2 \rangle 10. Q.E.D.
       PROOF: This is a contradiction.
```

**Definition 6.1.22** (Initial Segment). Let  $\leq$  be a partial order on **A** and  $t \in A$ . The *initial segment* up to t is the class

$$\operatorname{seg} t := \{ x \in \mathbf{A} \mid x < t \} .$$

**Definition 6.1.23** (Lexicographic Ordering). Let **R** be a partial order on **A** and **S** a partial order on **B**. The *lexicographic ordering*  $\leq$  on **A**  $\times$  **B** is defined by:

$$(a,b) \le (a',b') \Leftrightarrow (a\mathbf{R}a' \wedge a \ne a') \vee (a = a' \wedge b\mathbf{S}b')$$
.

**Proposition Schema 6.1.24** (Z). For any classes A, B, R and S, the following is a theorem:

If **R** is a partial order on **A** and **S** is a partial order on **B** then the lexicographic ordering on  $\mathbf{A} \times \mathbf{B}$  is a partial order.

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $\leq$  be the lexicographic ordering on  $\mathbf{A} \times \mathbf{B}$
- $\langle 1 \rangle 2. \leq \text{is reflexive.}$

PROOF: For any  $a \in \mathbf{A}$  and  $b \in \mathbf{B}$  we have a = a and  $b\mathbf{S}b$ , so  $(a, b) \leq (a, b)$ .

- $\langle 1 \rangle 3. \leq \text{is antisymmetric.}$ 
  - (2)1. Assume:  $(a,b) \le (a',b')$  and  $(a',b') \le (a,b)$
  - $\langle 2 \rangle 2$ .  $(a\mathbf{R}a' \wedge a \neq a') \vee (a = a' \wedge b\mathbf{S}b')$
  - $\langle 2 \rangle 3$ .  $(a' \mathbf{R} a \wedge a' \neq a) \vee (a' = a \wedge b \mathbf{S} b')$
  - $\langle 2 \rangle 4$ . Case: a = a'

PROOF: Then  $b\mathbf{S}b'$  and  $b'\mathbf{S}b$  hence b=b' and (a,b)=(a',b').

 $\langle 2 \rangle$ 5. Case:  $a \neq a'$ 

PROOF: Then  $a\mathbf{R}a'$  and  $a'\mathbf{R}a$  hence a=a' which is a contradiction.

- $\langle 1 \rangle 4$ .  $\leq$  is transitive.
  - $\langle 2 \rangle 1$ . Assume:  $(a_1, b_1) \le (a_2, b_2) \le (a_3, b_3)$
  - $\langle 2 \rangle 2$ .  $(a_1 \mathbf{R} a_2 \wedge a_1 \neq a_2) \vee (a_1 = a_2 \wedge b_1 \mathbf{S} b_2)$
  - $\langle 2 \rangle 3. \ (a_2 \mathbf{R} a_3 \wedge a_2 \neq a_3) \vee (a_2 = a_3 \wedge b_2 \mathbf{S} b_3)$
  - $\langle 2 \rangle 4$ . Case:  $a_1 \mathbf{R} a_2, a_1 \neq a_2, a_2 \mathbf{R} a_3, a_2 \neq a_3$ 
    - $\langle 3 \rangle 1. \ a_1 \mathbf{R} a_3$

PROOF: Since  $\mathbf{R}$  is transitive.

 $\langle 3 \rangle 2$ .  $a_1 \neq a_3$ 

PROOF: If  $a_1 = a_3$  then  $a_1 \mathbf{R} a_2$  and  $a_2 \mathbf{R} a_1$  so  $a_1 = a_2$  which is a contradiction.

 $\langle 2 \rangle 5$ . Case:  $a_1 \mathbf{R} a_2, a_1 \neq a_2, a_2 = a_3, b_2 \mathbf{S} b_3$ 

PROOF: Then  $a_1\mathbf{R}a_3$  and  $a_1 \neq a_3$ .

 $\langle 2 \rangle 6$ . Case:  $a_1 = a_2, b_1 \mathbf{S} b_2, a_2 \mathbf{R} a_3, a_2 \neq a_3$ 

PROOF: Then  $a_1 \mathbf{R} a_3$  and  $a_1 \neq a_3$ .

 $\langle 2 \rangle 7$ . Case:  $a_1 = a_2, b_1 \mathbf{S} b_2, a_2 = a_3, b_2 \mathbf{S} b_3$ 

PROOF: Then  $a_1 = a_3$  and  $b_1 \mathbf{S} b_3$ .

## 6.2 Linear Orders

**Definition 6.2.1** (Linear Ordering). Let **A** be a class. A *linear ordering* or *total ordering* on **A** is a partial ordering  $\leq$  on **A** that is *total*, i.e.

$$\forall x,y \in \mathbf{A}.x \leq y \vee y \leq x$$

We often use the symbol < for a linear ordering, and then write x < y for  $(x,y) \in <$ .

**Proposition Schema 6.2.2** (Trichotomy (Z)). For any classes **A** and  $\leq$ , the following is a theorem:

Assume  $\leq$  be a linear ordering on **A**. For any  $x, y \in \mathbf{A}$ , exactly one of x < y, x = y, y < x holds.

Proof: Immediate from definitions.  $\Box$ 

**Proposition Schema 6.2.3** (Z). For any classes A and <, the following is a theorem:

Let < be a transitive relation on  $\mathbf{A}$  that satisfies trichotomy. Define  $\leq$  on  $\mathbf{A}$  by  $x \leq y$  iff x < y or x = y. Then  $\leq$  is a linear ordering on  $\mathbf{A}$  and x < y iff  $x \leq y$  and  $x \neq y$ .

#### Proof:

 $\langle 1 \rangle 1$ . < is reflexive.

PROOF: By definition we have  $\forall x \in \mathbf{A}.x \leq x$ .

- $\langle 1 \rangle 2$ .  $\leq$  is antisymmetric.
  - $\langle 2 \rangle 1$ . Assume:  $x \leq y$  and  $y \leq x$
  - $\langle 2 \rangle 2$ . x < y or x = y
  - $\langle 2 \rangle 3$ . y < x or y = x
  - $\langle 2 \rangle 4$ . We cannot have x < y and y < x

PROOF: Trichotomy.

- $\langle 2 \rangle 5. \ x = y$
- $\langle 1 \rangle 3. \leq \text{is transitive.}$ 
  - $\langle 2 \rangle 1$ . Assume:  $x \leq y$  and  $y \leq z$
  - $\langle 2 \rangle 2$ . x < y or x = y
  - $\langle 2 \rangle 3$ . y < z or y = z
  - $\langle 2 \rangle 4$ . Case: x < y and y < z

PROOF: Then x < z by transitivity, so  $x \le z$ .

 $\langle 2 \rangle 5$ . Case: x = y

PROOF: Then we have  $y \leq z$  and so  $x \leq z$ .

 $\langle 2 \rangle 6$ . Case: y = z

PROOF: Then we have  $x \leq y$  and so  $x \leq z$ .

 $\langle 1 \rangle 4. \leq \text{is total.}$ 

PROOF: Immediate from trichotomy.

**Proposition Schema 6.2.4** (Z). For any classes **A** and **R**, the following is a theorem:

If  $\mathbf{R}$  is a linear ordering on  $\mathbf{A}$  then  $\mathbf{R}^{-1}$  is also a linear ordering on  $\mathbf{A}$ .

#### PROOF

 $\langle 1 \rangle 1$ .  $\mathbf{R}^{-1}$  is a partial order on  $\mathbf{A}$ .

Proof: Proposition 6.1.2.

 $\langle 1 \rangle 2$ .  $\mathbf{R}^{-1}$  is total.

```
\langle 2 \rangle 1. Let: x, y \in \mathbf{A}

\langle 2 \rangle 2. x \mathbf{R} y or y \mathbf{R} x.

\langle 2 \rangle 3. y \mathbf{R}^{-1} x or x \mathbf{R}^{-1} y.
```

**Proposition Schema 6.2.5** (Z). For any classes **A**, **B**, **F**, **R**, **S**, the following is a theorem:

Assume **R** is a linear order on **A**, **S** is a partial order on **B**, and **F** :  $\mathbf{A} \to \mathbf{B}$ . If **F** is strictly monotone then it is injective.

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $x, y \in \mathbf{A}$
- $\langle 1 \rangle 2$ . Assume:  $x \neq y$

PROVE:  $\mathbf{F}(x) \neq \mathbf{F}(y)$ 

 $\langle 1 \rangle 3$ . Assume: w.l.o.g.  $x \mathbf{R} y$ 

PROOF:  $\mathbf{R}$  is total.

 $\langle 1 \rangle 4$ .  $\mathbf{F}(x)\mathbf{SF}(y)$  and  $\mathbf{F}(x) \neq \mathbf{F}(y)$ 

PROOF: **F** is strictly monotone.

**Proposition Schema 6.2.6** (Z). For any classes A, B,  $\leq$ ,  $\preccurlyeq$  and F, the following is a theorem:

Assume  $\leq$  is a linear order on  $\mathbf{A}$  and  $\leq$  is a linear order on  $\mathbf{B}$ . Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{F}$  is strictly monotone. For all  $x, y \in \mathbf{A}$ , if  $\mathbf{F}(x) \prec \mathbf{F}(y)$  then x < y.

#### Proof:

- $\langle 1 \rangle 1$ .  $\mathbf{F}(x) \neq \mathbf{F}(y)$  and  $\mathbf{F}(y) \not\prec \mathbf{F}(x)$
- PROOF: Trichotomy.  $\langle 1 \rangle 2$ .  $x \neq y$  and  $y \not< x$

PROOF:  $\mathbf{F}$  is strictly monotone.

 $\langle 1 \rangle 3. \ x < y$ 

PROOF: Trichotomy.

**Corollary Schema 6.2.6.1** (Z). For any classes A, B,  $\leq$ ,  $\preccurlyeq$  and F, the following is a theorem:

Assume  $\leq$  is a linear order on  $\mathbf{A}$  and  $\preccurlyeq$  is a linear order on  $\mathbf{B}$ . Assume  $\mathbf{F}: \mathbf{A} \to \mathbf{B}$  and  $\mathbf{F}$  is strictly monotone. Then  $\mathbf{F}$  is an order isomorphism.

**Proposition Schema 6.2.7** (Z). For any classes **A**, **B**, **F** and **S**, the following is a theorem:

Assume **S** is a linear order on **B** and **F**:  $\mathbf{A} \rightarrow \mathbf{B}$ . Define **R** on **A** by  $x\mathbf{R}y$  if and only if  $\mathbf{F}(x)\mathbf{SF}(y)$ . Then **R** is a linear order on **A**.

#### Proof:

 $\langle 1 \rangle 1$ . **R** is a partial order on **A**.

Proof: Proposition 6.1.3.

```
\langle 1 \rangle 2. R is total.
    PROOF: For all x, y \in \mathbf{A} we have \mathbf{F}(x)\mathbf{SF}(y) or \mathbf{F}(y)\mathbf{SF}(x).
```

Corollary Schema 6.2.7.1 (Z). For any classes A, B and R, the following is a theorem:

Assume **R** be a linear order on **A** and **B**  $\subseteq$  **A**. Then **R**  $\cap$  **B**<sup>2</sup> is a linear order on **B**.

Proposition Schema 6.2.8 (Z). For any classes A, B, R and S, the following is a theorem:

Assume  $\mathbf{R}$  is a linear order on  $\mathbf{A}$  and  $\mathbf{S}$  is a linear order on  $\mathbf{B}$ . Then the lexicographic ordering is a linear order on  $\mathbf{A} \times \mathbf{B}$ .

```
Proof:
```

```
\langle 1 \rangle 1. Let: \leq be the lexicographic order on \mathbf{A} \times \mathbf{B}
\langle 1 \rangle 2. \leq is a partial order.
   Proof: Proposition 6.1.24.
\langle 1 \rangle 3. \leq \text{is total.}
   \langle 2 \rangle 1. Let: a, a' \in \mathbf{A} and b, b' \in \mathbf{B}
   \langle 2 \rangle 2. Case: a\mathbf{R}a' and a \neq a'
       PROOF: Then (a, b) \leq (a', b').
    \langle 2 \rangle 3. Case: a = a'
       PROOF: We have b\mathbf{S}b' or b'\mathbf{S}b, so (a,b) \leq (a',b') or (a',b') \leq (a,b).
   \langle 2 \rangle 4. Case: a' \mathbf{R} a and a \neq a'
       PROOF: Then (a', b') \leq (a, b).
```

#### 6.3 Well Orderings

**Definition 6.3.1** (Well Ordering). A well ordering on a class **A** is a wellfounded linear ordering on **A**.

**Proposition 6.3.2** (Z). Let S be a well ordering of the set B and  $f: A \to B$  a function. Define R on A by xRy if and only if F(x)SF(y). Then R well orders A.

```
\langle 1 \rangle 1. R linearly orders A.
   Proof: Proposition 6.2.7.
\langle 1 \rangle 2. Every nonempty subset of A has a least element.
   \langle 2 \rangle 1. Let: C be a nonempty subset of A.
   \langle 2 \rangle 2. Let: y be the least element of f(C).
   \langle 2 \rangle 3. PICK x \in C such that f(x) = y.
   \langle 2 \rangle 4. x is least in C.
```

**Proposition Schema 6.3.3** (Z). For any classes **A**, **B** and **R**, the following is a theorem:

Assume **R** well orders **B** and  $\mathbf{A} \subseteq \mathbf{B}$ . Then  $\mathbf{R} \cap \mathbf{A}^2$  well orders **A**.

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $\mathbf{R}' = \mathbf{R} \cap \mathbf{A}^2$
- $\langle 1 \rangle 2$ . **R'** linearly orders **A**.

Proof: Corollary 6.2.7.1.

 $\langle 1 \rangle 3$ . **R**' is well founded.

Proof: Proposition 2.4.4.

**Proposition Schema 6.3.4** (ZFC). For any classes **A**, **B**, **F** and **S**, the following is a theorem:

Assume **S** well orders **B** and **F** :  $\mathbf{A} \rightarrow \mathbf{B}$ . Define **R** on **A** by  $x\mathbf{R}y$  if and only if  $\mathbf{F}(x)\mathbf{SF}(y)$ . Then **R** well orders **A**.

#### Proof:

 $\langle 1 \rangle 1$ . **R** linearly orders **A**.

Proof: Proposition 6.2.7.

- $\langle 1 \rangle 2$ . For all  $t \in \mathbf{A}$  we have  $\{x \in \mathbf{A} \mid x\mathbf{R}t \land x \neq t\}$  is a set.
  - $\langle 2 \rangle 1$ . Let:  $t \in \mathbf{A}$
  - $\langle 2 \rangle 2$ . Let:  $S = \{ y \in \mathbf{B} \mid y\mathbf{SF}(t) \land y \neq \mathbf{F}(t) \}$
  - $\langle 2 \rangle 3$ . Let: P(x,y) be the property  $\mathbf{F}(y) = x$
  - $\langle 2 \rangle 4$ . For all  $x \in S$  there exists at most one y such that P(x, y) PROOF: **F** is injective.
  - $\langle 2 \rangle$ 5. Let:  $T = \{ y \mid \exists x \in S.P(x,y) \}$

PROOF: Axiom of Replacement.

- $\langle 2 \rangle 6. \ T = \{ x \in \mathbf{A} \mid x\mathbf{R}t \land x \neq t \}$
- $\langle 1 \rangle 3$ . Every nonempty subset of **A** has a least element.
  - $\langle 2 \rangle 1$ . Let: S be a nonempty subset of **A**.
  - $\langle 2 \rangle 2$ . **F**(S) is a nonempty subset of **B**

PROOF: Axiom of Replacement.

- $\langle 2 \rangle 3$ . Let: y be the least element of  $\mathbf{F}(S)$ .
- $\langle 2 \rangle 4$ . PICK  $x \in S$  such that  $\mathbf{F}(x) = y$ .
- $\langle 2 \rangle 5$ . x is least in S.

П

**Proposition 6.3.5** (Z). For any well ordered sets A and B, the lexicographic order well orders  $A \times B$ .

#### Proof:

 $\langle 1 \rangle 1$ .  $A \times B$  is linearly ordered.

Proof: Proposition 6.2.8.

- $\langle 1 \rangle 2$ . Every nonempty subset of  $A \times B$  has a least element.
  - $\langle 2 \rangle 1$ . Let: S be a nonempty subset of  $A \times B$ .
  - $\langle 2 \rangle 2$ . Let: a be the least element of  $\{x \in A \mid \exists y \in B.(x,y) \in S\}$ .
  - $\langle 2 \rangle 3$ . Let: b be the least element of  $\{ y \in B \mid (a, y) \in S \}$ .

(2)4. (a,b) is least in S.

**Definition 6.3.6** (End Extension). Let A and B be well ordered sets. Then B is an *end extension* of A iff  $A \subseteq B$  and:

- Whenever  $x, y \in A$  then  $x \leq_A y$  iff  $x \leq_B y$ .
- Whenever  $x \in A$  and  $y \in B A$  then x < y.

**Theorem 6.3.7** (Z). Let  $\leq$  be a linear ordering on A. Assume that, for any  $B \subseteq A$  such that  $\forall t \in A$ . seg  $t \subseteq B \Rightarrow t \in B$ , we have B = A. Then  $\leq$  is a well ordering on A.

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $C \subseteq A$  be nonempty.
- $\langle 1 \rangle 2$ . Let:  $B = \{ t \in A \mid \forall x \in C.t < x \}$
- $\langle 1 \rangle 3. \ B \cap C = \emptyset$
- $\langle 1 \rangle 4$ .  $B \neq A$
- $\langle 1 \rangle$ 5. PICK  $t \in A$  such that  $\operatorname{seg} t \subseteq B$  and  $t \notin B$
- $\langle 1 \rangle 6$ . t is least in C.

**Proposition Schema 6.3.8** (Z). For any classes A, B, F, G,  $\leq$  and  $\preccurlyeq$ , the following is a theorem:

Assume  $\leq$  well orders  $\mathbf{A}$  and  $\leq$  well orders  $\mathbf{B}$ . Assume  $\mathbf{F}$  and  $\mathbf{G}$  are order isomorphisms between  $\mathbf{A}$  and  $\mathbf{B}$ . Then  $\mathbf{F} = \mathbf{G}$ .

### Proof:

- $\langle 1 \rangle 1$ . For all  $x \in \mathbf{A}$ , if  $\forall t < x.\mathbf{F}(t) = \mathbf{G}(t)$ , then  $\mathbf{F}(x) = \mathbf{G}(x)$ 
  - $\langle 2 \rangle 1$ . Let:  $x \in \mathbf{A}$
  - $\langle 2 \rangle 2$ . Assume:  $\forall t < x. \mathbf{F}(t) = \mathbf{G}(t)$
  - $\langle 2 \rangle 3$ .  $\mathbf{F}(\operatorname{seg} x) = \mathbf{G}(\operatorname{seg} x)$
  - $\langle 2 \rangle 4$ .  $\mathbf{F}(x)$  is the least element of  $\mathbf{B} \mathbf{F}(\operatorname{seg} x)$
  - $\langle 2 \rangle 5$ .  $\mathbf{G}(x)$  is the least element of  $\mathbf{B} \mathbf{G}(\operatorname{seg} x)$
  - $\langle 2 \rangle 6. \ \mathbf{F}(x) = \mathbf{G}(x)$
- $\langle 1 \rangle 2. \ \forall x \in \mathbf{A}.\mathbf{F}(x) = \mathbf{G}(x)$

PROOF: Transfinite induction.

**Theorem 6.3.9** (ZFC). Let A and B be well ordered sets. Then one of the following holds:  $A \cong B$ ; there exists  $b \in B$  such that  $A \cong \operatorname{seg} b$ ; there exists  $a \in A$  such that  $\operatorname{seg} a \cong B$ .

- $\langle 1 \rangle 1$ . PICK e that is not in A or B.
- $\langle 1 \rangle$ 2. Let:  $F: A \to B \cup \{e\}$  be the function defined by transfinite recursion thus:

$$F(t) = \begin{cases} \text{the least element of } B - F(\sec t) & \text{if } B - F(\sec t) \neq \emptyset \\ e & \text{if } B - F(\sec t) = \emptyset \end{cases}$$

```
\begin{split} &\langle 1 \rangle 3. \text{ Case: } e \in \operatorname{ran} F \\ &\langle 2 \rangle 1. \text{ Let: } t \text{ be least such that } F(t) = e \\ &\langle 2 \rangle 2. F \upharpoonright \operatorname{seg} t : \operatorname{seg} t \cong B \\ &\langle 1 \rangle 4. \text{ Case: } \operatorname{ran} F = B \\ &\operatorname{PROOF: We have } F: A \cong B \\ &\langle 1 \rangle 5. \text{ Case: } \operatorname{ran} F \subsetneq B \\ &\langle 2 \rangle 1. \text{ Let: } b \text{ be the least element of } B - \operatorname{ran} F \\ &\langle 2 \rangle 2. F: A \cong \operatorname{seg} b \end{split}
```

## Chapter 7

## Ordinal Numbers

## 7.1 Ordinals

**Definition 7.1.1** (Ordinal). An *ordinal (number)* is a transitive set  $\alpha$  such that, for all  $x, y \in \alpha$ , we have  $x \in y$  or x = y or  $y \in x$ .

Given ordinals  $\alpha$ ,  $\beta$ , we write  $\alpha < \beta$  iff  $\alpha \in \beta$ , and  $\alpha \le \beta$  iff  $\alpha \in \beta$  or  $\alpha = \beta$ . Let **On** be the class of ordinals.

**Proposition 7.1.2** (S). If  $\beta$  is an ordinal and  $\alpha \in \beta$  then  $\alpha$  is an ordinal.

```
Proof:
\langle 1 \rangle 1. Let: \beta be an ordinal.
\langle 1 \rangle 2. Let: \alpha \in \beta
\langle 1 \rangle 3. \alpha is a transitive set.
   \langle 2 \rangle 1. Let: x \in y \in \alpha
   \langle 2 \rangle 2. \ y \in \beta
       PROOF: \beta is a transitive set.
   \langle 2 \rangle 3. \ x \in \beta
       PROOF: \beta is a transitive set.
    \langle 2 \rangle 4. x \in \alpha or x = \alpha or \alpha \in x
   \langle 2 \rangle 5. \ x \neq \alpha
       PROOF: We cannot have x \in y \in x by the Axiom of Foundation.
       PROOF: We cannot have x \in y \in \alpha \in x by the Axiom of Foundation.
    \langle 2 \rangle 7. \ x \in \alpha
\langle 1 \rangle 4. \ \forall x, y \in \alpha. x \in y \lor x = y \lor y \in x
   PROOF: This holds because if x, y \in \alpha then x, y \in \beta since \beta is a transitive
   set.
Proposition 7.1.3 (S). \emptyset is an ordinal.
```

**Definition 7.1.4** (Successor). For any ordinal  $\alpha$ , the *successor* of  $\alpha$  is

$$\alpha^+ := \alpha \cup \{\alpha\} .$$

**Proposition 7.1.5** (S). For any ordinal  $\alpha$ , the successor  $\alpha^+$  is an ordinal.

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $\alpha$  be an ordinal.
- $\langle 1 \rangle 2$ .  $\alpha^+$  is a transitive set.
  - $\langle 2 \rangle 1$ . Let:  $x \in y \in \alpha^+$
  - $\langle 2 \rangle 2$ .  $y \in \alpha$  or  $y = \alpha$
  - $\langle 2 \rangle 3. \ x \in \alpha$

PROOF: If  $y \in \alpha$  then this follows because  $\alpha$  is a transitive set. If  $y = \alpha$  this follows immediately.

- $\langle 2 \rangle 4. \ x \in \alpha^+$
- $\langle 1 \rangle 3$ . For any  $x, y \in \alpha^+$ , we have  $x \in y$  or x = y or  $y \in X$ .
  - $\langle 2 \rangle 1$ . Let:  $x, y \in \alpha^+$
  - $\langle 2 \rangle 2$ . Case:  $x \in \alpha$  and  $y \in \alpha$

PROOF: Then  $x \in y$  or x = y or  $y \in x$  because  $\alpha$  is an ordinal.

 $\langle 2 \rangle 3$ . Case:  $x \in \alpha$  and  $y = \alpha$ 

PROOF: Then  $x \in y$ .

 $\langle 2 \rangle 4$ . Case:  $x = \alpha$  and  $y \in \alpha$ 

PROOF: Then  $y \in x$ .

 $\langle 2 \rangle$ 5. Case:  $x = \alpha$  and  $y = \alpha$ 

PROOF: Then x = y.

П

**Proposition 7.1.6** (S). For ordinals  $\alpha$  and  $\beta$ , we have  $\alpha \subseteq \beta$  if and only if  $\alpha \leq \beta$ .

#### Proof:

 $\langle 1 \rangle 1$ . If  $\alpha \leq \beta$  then  $\alpha \subseteq \beta$ .

PROOF: Since  $\beta$  is a transitive set.

- $\langle 1 \rangle 2$ . If  $\alpha \subseteq \beta$  then  $\alpha \leq \beta$ .
  - $\langle 2 \rangle 1$ . Assume:  $\alpha \subsetneq \beta$

Prove:  $\alpha \in \beta$ 

 $\langle 2 \rangle 2$ . Pick  $u \in \beta - \alpha$  such that  $u \cap (\beta - \alpha) = \emptyset$ 

PROOF: Axiom of Foundation.

- $\langle 2 \rangle 3. \ u \in \beta$
- $\langle 2 \rangle 4. \ u \notin \alpha$
- $\langle 2 \rangle$ 5. For all  $v \in \beta$ , if  $v \notin \alpha$  then  $v \notin u$

PROOF: From  $\langle 2 \rangle 2$ 

- $\langle 2 \rangle 6$ . For all  $v \in u$  we have  $v \in \alpha$
- $\langle 2 \rangle 7$ . For all  $w \in \alpha$  we have  $w \in u$ 
  - $\langle 3 \rangle 1$ . Let:  $w \in \alpha$
  - $\langle 3 \rangle 2. \ w \in \beta$

Proof:  $\langle 2 \rangle 1$ 

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\langle 3 \rangle 3. \ w \in u \text{ or } w = u \text{ or } u \in w
            PROOF: \beta is an ordinal, \langle 2 \rangle 3, \langle 3 \rangle 2
        \langle 3 \rangle 4. \ w \neq u
            Proof: \langle 2 \rangle 4, \langle 3 \rangle 1
        \langle 3 \rangle 5. \ u \notin w
            PROOF: \langle 2 \rangle 4, \langle 3 \rangle 1, \alpha is a transitive set.
        \langle 3 \rangle 6. \ w \in u
            Proof: \langle 3 \rangle 3, \langle 3 \rangle 4, \langle 3 \rangle 5
    \langle 2 \rangle 8. \ u = \alpha
        PROOF: \langle 2 \rangle 6, \langle 2 \rangle 7, Axiom of Extensionality.
    \langle 2 \rangle 9. \ \alpha \in \beta
        Proof: \langle 2 \rangle 3, \langle 2 \rangle 8
Proposition 7.1.7 (S). The intersection of any nonempty set of ordinals is an
ordinal.
Proof:
\langle 1 \rangle 1. Let: x be a nonempty set of ordinals.
\langle 1 \rangle 2. \bigcap x is a transitive set.
    Proof: Proposition 1.6.6.
\langle 1 \rangle 3. \ \forall y, z \in \bigcap x. y \in z \lor y = z \lor z \in y
    \langle 2 \rangle 1. Pick \alpha \in x
    \langle 2 \rangle 2. \ y, z \in \alpha
    \langle 2 \rangle 3. \ y \in z \text{ or } y = z \text{ or } z \in y
Theorem 7.1.8 (S). For any ordinals \alpha and \beta, we have \alpha < \beta or \alpha = \beta or
\beta < \alpha.
Proof:
\langle 1 \rangle 1. Let: \alpha and \beta be ordinals.
\langle 1 \rangle 2. Case: \alpha \cap \beta = \alpha or \alpha \cap \beta = \beta
    \langle 2 \rangle 1. Assume: for a contradiction \alpha \cap \beta \subseteq \alpha and \alpha \cap \beta \subseteq \beta
    \langle 2 \rangle 2. \alpha \cap \beta \in \alpha and \alpha \cap \beta \in \beta
        Proof: Propositions 7.1.6, 7.1.7.
    \langle 2 \rangle 3. \ \alpha \cap \beta \in \alpha \cap \beta
    \langle 2 \rangle 4. Q.E.D.
        PROOF: This contradicts the Axiom of Foundation.
\langle 1 \rangle 3. \ \alpha \subseteq \beta \text{ or } \beta \subseteq \alpha
\langle 1 \rangle 4. \alpha \in \beta or \alpha = \beta or \beta \in \alpha
    Proof: Proposition 7.1.6.
Theorem 7.1.9 (S). The class On is well ordered by \leq.
Proof:
\langle 1 \rangle 1. For every ordinal \alpha we have \alpha \leq \alpha.
```

```
PROOF: Immediate from definition. 
 \langle 1 \rangle 2. For any ordinals \alpha, \beta and \gamma, if \alpha \leq \beta \leq \gamma then \alpha \leq \gamma. 
 PROOF: Since \gamma is a transitive set. 
 \langle 1 \rangle 3. For any ordinals \alpha and \beta, if \alpha \leq \beta \leq \alpha then \alpha = \beta. 
 PROOF: If \alpha \neq \beta then \alpha \in \beta \in \alpha contradicting the Axiom of Foundation. 
 \langle 1 \rangle 4. For any ordinals \alpha and \beta, either \alpha \leq \beta or \beta \leq \alpha. 
 PROOF: Theorem 7.1.8. 
 \langle 1 \rangle 5. Any set of ordinals x has a least element. 
 PROOF: Axiom of Foundation.
```

**Proposition 7.1.10.** For any nonempty set of ordinals x, we have  $\bigcap x$  is the infimum of x in  $\mathbf{On}$ .

PROOF: It is an ordinal by Proposition 7.1.7. The fact it is the infimum is immediate from definitions.  $\Box$ 

**Proposition 7.1.11.** For any nonempty set of ordinals x, we have  $\bigcup x$  is the supremum of x in On.

```
Proof:
```

- $\langle 1 \rangle 1$ .  $\bigcup x$  is an ordinal.
  - $\langle 2 \rangle 1$ .  $\bigcup x$  is a transitive set.

Proof: Proposition 1.6.3.

- $\langle 2 \rangle 2$ . For all  $y, z \in \bigcup x$  we have  $y \in z$  or y = z or  $z \in y$ 
  - $\langle 3 \rangle 1$ . PICK  $\alpha, \beta \in x$  such that  $y \in \alpha$  and  $z \in \beta$ .
  - $\langle 3 \rangle 2$ . Assume: w.l.o.g.  $\alpha \leq \beta$ .
  - $\langle 3 \rangle 3. \ y, z \in \beta$
- $\langle 3 \rangle 4. \ y \in z \text{ or } y = z \text{ or } z \in y.$

**Definition 7.1.12** (Successor Ordinal). An ordinal  $\alpha$  is a *successor ordinal* iff there exists an ordinal  $\beta$  such that  $\alpha = \beta^+$ .

**Definition 7.1.13** (Limit Ordinal). An ordinal is a *limit ordinal* iff it is neither 0 nor a successor ordinal.

## 7.2 Natural Numbers

**Definition 7.2.1** (Natural Number). An ordinal n is a *natural number* iff either it is 0, or it is a successor ordinal and every element of n is either 0 or a successor ordinal.

We write  $\omega$  or N for the class of natural numbers.

**Theorem Schema 7.2.2** (Mathematical Induction (S)). For any predicate P(x), the following is a theorem:

Assume P(0). Assume that, for any ordinal  $\alpha$ , if  $P(\alpha)$  then  $P(\alpha^+)$ . Then, for every natural number n, we have P(n).

```
Proof:
\langle 1 \rangle 1. Assume: n is a natural number such that \neg P(n)
\langle 1 \rangle 2. Let: y = \{ \gamma \in n^+ \mid \neg P(\gamma) \}
\langle 1 \rangle 3. \ y \neq \emptyset
\langle 1 \rangle 4. Pick \gamma_0 \in y such that y \cap \gamma_0 = \emptyset
\langle 1 \rangle 5. Case: \gamma_0 = 0
   PROOF: Then \neg P(0)
\langle 1 \rangle 6. Case: \gamma_0 \neq 0
   \langle 2 \rangle 1. PICK \gamma_1 such that \gamma_0 = \gamma_1^+
       PROOF: Either \gamma_0 = n or \gamma_0 \in n, and in either case \gamma_0 is a successor because
       n is a natural number.
   \langle 2 \rangle 2. \ \gamma_1 \notin y
       PROOF: From \langle 1 \rangle 4 since \gamma_1 \in \gamma_0.
   \langle 2 \rangle 3. \ \gamma_1 \in n^+
       PROOF: Since \gamma_1 \in \gamma_0 \in n^+ and n^+ is a transitive set.
   \langle 2 \rangle 4. P(\gamma_1) and \neg P(\gamma_1^+)
       Proof: \langle 1 \rangle 2, \langle 2 \rangle 2, \langle 2 \rangle 3
Theorem 7.2.3 (Z). The class of natural numbers \omega is a set.
\langle 1 \rangle 1. PICK a set I such that \emptyset \in I and \forall x \in I.x \cup \{x\} \in I.
   Proof: Axiom of Infinity.
\langle 1 \rangle 2. Every natural number is in I.
   PROOF: By induction.
\langle 1 \rangle 3. Q.E.D.
   Proof: By an Axiom of Comprehension.
Theorem 7.2.4. \omega is a limit ordinal.
PROOF:
\langle 1 \rangle 1. \omega is a transitive set.
   \langle 2 \rangle 1. Let: m \in n \in \omega
            Prove: m is a natural number.
   \langle 2 \rangle 2. Assume: m \neq 0
   \langle 2 \rangle 3. m is a successor ordinal.
       PROOF: Since m \in n and n is a natural number.
   \langle 2 \rangle 4. Every element of m is either 0 or a successor ordinal.
       \langle 3 \rangle 1. Let: x \in m
       \langle 3 \rangle 2. \ x \in n
          PROOF: Since n is a transitive set.
       \langle 3 \rangle 3. x = 0 or x is a successor ordinal.
\langle 1 \rangle 2. For all m, n \in \omega, either m \in n or m = n or n \in m.
   PROOF: Theorem 7.1.8.
\langle 1 \rangle 3. \ \omega \neq 0
```

PROOF: Since  $0 \in \omega$ .

- $\langle 1 \rangle 4$ .  $\omega$  is not a successor ordinal.
  - $\langle 2 \rangle 1$ . Assume: for a contradiction  $\omega = n^+$
  - $\langle 2 \rangle 2$ . *n* is a natural number.
  - $\langle 2 \rangle 3$ .  $\omega$  is a natural number.
  - $\langle 2 \rangle 4. \ \omega \in \omega$
  - $\langle 2 \rangle$ 5. Q.E.D.

PROOF: This contradicts the Axiom of Foundation.

## 7.3 Transitive Closure of a Set

**Definition 7.3.1** (Transitive Closure). Let x be a set. The *transitive closure* of x is the class

$$TC(x) = \bigcap \{y \mid y \text{ is a transitive set } \land x \in y\}$$
.

**Theorem 7.3.2** (ZF). For any set x, the transitive closure TC(x) is a set.

Proof:

- $\langle 1 \rangle 1$ . Let: E(n,y,x) be the predicate: There exists a function  $f:n^+ \to \mathbf{V}$  such that f(0)=y and f(n)=x and  $\forall m \in n. f(m) \in f(m^+)$
- $\langle 1 \rangle 2. \ \forall x, y. E(0, y, x) \Leftrightarrow y = x$
- $\langle 1 \rangle 3$ . Let: x be a set.
- $\langle 1 \rangle 4$ . Let:  $\mathbf{W} = \{ y \mid \exists n \in \mathbb{N}. E(n, y, x) \}$
- $\langle 1 \rangle 5$ . **W** is a set.
  - $\langle 2 \rangle 1$ . For  $n \in \mathbb{N}$ ,

Let: 
$$\mathbf{W}_n = \{y \mid E(n, y, x)\}$$

- $\langle 2 \rangle 2$ . For all  $n \in \mathbb{N}$ ,  $\mathbf{W}_n$  is a set.
  - $\langle 3 \rangle 1$ .  $\mathbf{W}_0$  is a set.

PROOF:  $\mathbf{W}_0 = \{x\}.$ 

- $\langle 3 \rangle 2$ . For all  $n \in \mathbb{N}$ , if  $\mathbf{W}_n$  is a set then  $\mathbf{W}_{n^+}$  is a set.
  - $\langle 4 \rangle 1$ .  $\mathbf{W}_{n^+} = \bigcup \mathbf{W}_n$ 
    - $\langle 5 \rangle 1. \ \mathbf{W}_{n^+} \subseteq \bigcup \mathbf{W}_n$ 
      - $\langle 6 \rangle 1$ . Let:  $y \in \mathbf{W}_{n^+}$
      - $\langle 6 \rangle 2$ .  $E(n^+, y, x)$
      - (6)3. PICK  $f: n^{++} \to \mathbf{V}$  such that f(0) = y and  $f(n^+) = x$  and  $\forall m \in n^+. f(m) \in f(m^+)$
      - $\langle 6 \rangle 4. \ y \in f(1) \text{ and } E(n, f(1), x)$

PROOF: Define  $g: n^+ \to \mathbf{V}$  by  $g(m) = f(m^+)$ . Then g(0) = f(1) and g(n) = x and  $\forall m \in n. g(m) \in g(m^+)$ .

- $\langle 6 \rangle 5. \ y \in f(1) \in \mathbf{W}_n$
- $\langle 5 \rangle 2. \bigcup \mathbf{W}_n \subseteq \mathbf{W}_{n^+}$ 
  - $\langle 6 \rangle 1$ . Let:  $y \in z \in \mathbf{W}_n$
  - $\langle 6 \rangle 2$ . PICK  $f: n^+ \to \mathbf{V}$  such that f(0) = z and f(n) = x and  $\forall m \in n. f(m) \in f(m^+)$

```
\langle 6 \rangle 3. Let: g: n^{++} \to \mathbf{V} be the function defined by g(0) = y,
                                      g(m^+) = f(m)
                   (6)4. \ g(0) = y \text{ and } g(n^+) = x \text{ and } \forall m \in n^+. f(m) \in f(m^+)
                   \langle 6 \rangle 5. E(n^+, y, x)
   \langle 2 \rangle 3. \ \mathbf{W} = \bigcup \{ \mathbf{W}_n \mid n \in \mathbb{N} \}
   \langle 2 \rangle 4. \{ \mathbf{W}_n \mid n \in \mathbb{N} \} is a set.
       Proof: Axiom of Replacement.
   \langle 2 \rangle 5. W is a set.
       PROOF: Axiom of Union.
\langle 1 \rangle 6. W is a transitive set.
   \langle 2 \rangle 1. Let: y \in z \in \mathbf{W}
   \langle 2 \rangle 2. Pick n \in \mathbb{N} and f: n^+ \to \mathbf{V} such that f(0) = z, f(n) = x and
            \forall m \in n. f(m^+) \in f(m)
   \langle 2 \rangle 3. Let: g: n^{++} \to \mathbf{V} be the function with g(0) = y and g(m^+) = f(m)
   \langle 2 \rangle 4. \ \ g(0) = y \ \text{and} \ \ g(n^+) = x \ \text{and} \ \ \forall m \in n^{++}.g(m^+) \in g(m)
   \langle 2 \rangle 5. \ y \in \mathbf{W}
\langle 1 \rangle 7. \ x \in \mathbf{W}
   \langle 2 \rangle 1. Let: f: 0^+ \to \mathbf{V} be the function f = \{(0, x)\}
   \langle 2 \rangle 2. f(0) = x and f(0) = x and vacuously \forall m \in 0. f(m) \in f(m^+)
   \langle 2 \rangle 3. E(0,x,x)
\langle 1 \rangle 8. For any set w', if w' is a transitive set that contains x then \mathbf{W} \subseteq w'
   \langle 2 \rangle 1. Let: w' be a transitive set that contains x
   \langle 2 \rangle 2. \ \forall n \in \mathbb{N}. \forall y. E(n, y, x) \Rightarrow y \in w'
       \langle 3 \rangle 1. \ \forall y. E(0, y, x) \Rightarrow y \in w'
           \langle 4 \rangle 1. Let: y be a set.
           \langle 4 \rangle 2. Assume: E(0, y, x)
           \langle 4 \rangle 3. PICK a function f: 0^+ \to \mathbf{V} such that f(0) = y and f(0) = x and
                    \forall m \in 0. f(m) \in f(m^+)
           \langle 4 \rangle 4. x = y
              PROOF: x = f(0) = y
           \langle 4 \rangle 5. \ y \in w'
              Proof: \langle 2 \rangle 1
       \langle 3 \rangle 2. For all n \in \mathbb{N}, if \forall y. E(n, y, x) \Rightarrow y \in w', then \forall y. E(n^+, y, x) \Rightarrow y \in w'
           \langle 4 \rangle 1. Let: n \in \mathbb{N}
           \langle 4 \rangle 2. Assume: \forall y. E(n, y, x) \Rightarrow y \in w'
           \langle 4 \rangle 3. Let: y be a set.
           \langle 4 \rangle 4. Assume: E(n^+, y, x)
           \langle 4 \rangle5. PICK f: n^{++} \rightarrow \mathbf{V} such that f(0) = y and f(n^{+}) = x and
                    \forall m \in n^+.f(m) \in f(m^+)
           \langle 4 \rangle 6. Let: g: n^+ \to \mathbf{V} be the function g(m) = f(m^+)
           \langle 4 \rangle 7. \ \ g(0) = f(0^+) \ \text{and} \ \ g(n) = x \ \text{and} \ \forall m \in n. \ g(m) \in g(m^+)
           \langle 4 \rangle 8. \ E(n, f(0^+), x)
           \langle 4 \rangle 9. \ f(0^+) \in w'
              Proof: \langle 4 \rangle 5
           \langle 4 \rangle 10. \ y \in f(0^+)
```

```
PROOF: y = f(0) \in f(0^+) by \langle 4 \rangle 5. \langle 4 \rangle 11. y \in w' PROOF: From \langle 3 \rangle 2 and \langle 4 \rangle 10 since w' is transitive (\langle 2 \rangle 1). \langle 2 \rangle 3. \mathbf{W} \subseteq w'
```

**Proposition 7.3.3** (Z). For any ordinal numbers  $\alpha$  and  $\beta$ , if  $\alpha \cong \beta$  then  $\alpha = \beta$ .

```
Proof:
```

```
\langle 1 \rangle 1. Let: f : \alpha \cong \beta
\langle 1 \rangle 2. For all x \in \alpha, if \forall t < x. f(t) = t then f(x) = x
    \langle 2 \rangle 1. \ f(x) \subseteq x
        \langle 3 \rangle 1. Let: y \in f(x)
        \langle 3 \rangle 2. \ y \in \beta
        \langle 3 \rangle 3. Pick t \in \alpha such that f(t) = y
            Proof: f is surjective.
        \langle 3 \rangle 4. \ f(t) \in f(x)
        \langle 3 \rangle 5. \ t \in x
            PROOF: Since f is an order isomorphism.
        \langle 3 \rangle 6. f(t) = t
            Proof: Induction hypothesis.
        \langle 3 \rangle 7. \ y = t
        \langle 3 \rangle 8. \ y \in x
    \langle 2 \rangle 2. x \subseteq f(x)
        \langle 3 \rangle 1. Let: t \in x
        \langle 3 \rangle 2. \ f(t) \in f(x)
        \langle 3 \rangle 3. \ f(t) = t
        \langle 3 \rangle 4. \ t \in f(x)
\langle 1 \rangle 3. \ \forall x \in \alpha. f(x) = x
   PROOF: Transfinite induction.
\langle 1 \rangle 4. \ \alpha = \beta
   PROOF: Since \beta = \{f(t) \mid t \in \alpha\} = \{t \mid t \in \alpha\} = \alpha.
```

**Theorem 7.3.4** (ZFC). Every well-ordered set is isomorphic to a unique ordinal.

- $\langle 1 \rangle 1$ . For any well-ordered set A, there exists an ordinal  $\alpha$  such that  $A \cong \alpha$ .
  - $\langle 2 \rangle 1$ . Let: A be a well-ordered set.
  - $\langle 2 \rangle 2$ . Define the function E on A by transfinite recursion thus:

$$E(t) = \{ E(x) \mid x < t \}$$
  $(t \in A)$ .

- $\langle 2 \rangle 3$ . Let:  $\alpha = \{ E(x) \mid x \in A \}$
- $\langle 2 \rangle 4$ .  $\alpha$  is an ordinal.
  - $\langle 3 \rangle 1$ .  $\alpha$  is a transitive set.
    - $\langle 4 \rangle 1$ . Let:  $x \in y \in \alpha$
    - $\langle 4 \rangle 2$ . Pick  $t \in A$  such that y = E(t)

```
\langle 4 \rangle 3. \ x \in E(t) = \{ E(s) \mid s < t \}
          \langle 4 \rangle 4. Pick s < t such that x = E(s)
           \langle 4 \rangle 5. \ x \in \alpha
      \langle 3 \rangle 2. \alpha is well-ordered by \in.
          \langle 4 \rangle 1. Let: \langle = \{ (x, y) \in \alpha \mid x \in y \}
          \langle 4 \rangle 2. < is transitive.
              \langle 5 \rangle 1. Let: x, y, z \in \alpha with x \in y \in z
              \langle 5 \rangle 2. Pick t \in A such that z = E(t)
              \langle 5 \rangle 3. Pick s \in A such that s < t and y = E(s)
              \langle 5 \rangle 4. PICK r \in A such that r < s and x = E(r)
              \langle 5 \rangle 5. r < t
              \langle 5 \rangle 6. \ x \in z
          \langle 4 \rangle 3. < satisfies trichotomy.
              \langle 5 \rangle 1. Let: x, y \in \alpha
              \langle 5 \rangle 2. PICK s, t \in A such that E(s) = x and E(t) = y
              \langle 5 \rangle 3. Exactly one of s < t, s = t, t < s holds.
              \langle 5 \rangle 4. Case: s < t
                  \langle 6 \rangle 1. \ x \in y
                  \langle 6 \rangle 2. x \neq y and y \notin x
                     PROOF: Axiom of Regularity.
              \langle 5 \rangle 5. Case: s = t
                  \langle 6 \rangle 1. \ x = y
                  \langle 6 \rangle 2. x \notin y and y \notin x
                     Proof: Axiom of Regularity.
              \langle 5 \rangle 6. Case: t < s
                  Proof: Similar to \langle 5 \rangle 4.
          \langle 4 \rangle 4. \leq is a linear order on \alpha.
              Proof: Proposition 6.2.3.
          \langle 4 \rangle5. Every nonempty subset of \alpha has a least element.
              \langle 5 \rangle 1. Let: S be a nonempty subset of \alpha
              \langle 5 \rangle 2. Let: T = \{ x \in A \mid E(x) \in S \}
              \langle 5 \rangle 3. Let: t be the least element of T.
                       PROVE: E(t) is least in S
              \langle 5 \rangle 4. Let: y \in S
              \langle 5 \rangle 5. PICK s \in T such that E(s) = y
              \langle 5 \rangle 6. \ t \leq s
              \langle 5 \rangle 7. \ x \leq y
   \langle 2 \rangle5. E is surjective.
      PROOF: By definition of \alpha.
   \langle 2 \rangle 6. E is strictly monotone.
      PROOF: If s < t then E(s) \in E(t) by definition of E(t).
   \langle 2 \rangle7. Q.E.D.
      Proof: Corollary 6.2.6.1.
\langle 1 \rangle 2. For any ordinals \alpha and \beta, if \alpha \cong \beta then \alpha = \beta.
   Proof: Proposition 7.3.3.
```

Proof:

**Proposition 7.3.5** (Z). The class **On** is a transitive class. That is, every element of an ordinal is an ordinal.

```
\langle 1 \rangle 1. Let: \alpha be an ordinal.
\langle 1 \rangle 2. Let: \beta \in \alpha
\langle 1 \rangle 3. \beta is a transitive set.
    \langle 2 \rangle 1. Let: x \in y \in \beta
    \langle 2 \rangle 2. \ y \in \alpha
       PROOF: \alpha is transitive.
    \langle 2 \rangle 3. \ x \in \alpha
        Proof: \alpha is transitive.
    \langle 2 \rangle 4. \ x \in \beta
        PROOF: Since \{(x,y) \in \alpha^2 \mid x \in y\} is transitive.
\langle 1 \rangle 4. \beta is well ordered by \in.
    Proof: By Proposition 6.3.3.
Proposition 7.3.6 (ZFC). Given two ordinal numbers \alpha, \beta, exactly one of
\alpha \in \beta, \alpha = \beta, \beta \in \alpha holds.
Proof:
\langle 1 \rangle 1. At most one holds.
    PROOF: Since every ordinal is a transitive set and we never have \alpha \in \alpha.
\langle 1 \rangle 2. At least one holds.
    \langle 2 \rangle 1. Either \alpha \cong \beta or \exists t \in \beta . \alpha \cong \text{seg } t or \exists t \in \alpha . \text{seg } t \cong \beta .
    \langle 2 \rangle 2. Case: \alpha \cong \beta
        PROOF: Then \alpha = \beta by Proposition 7.3.3.
    \langle 2 \rangle 3. Case: There exists t \in \beta such that \alpha \cong \operatorname{seg} t
        \langle 3 \rangle 1. t is an ordinal number.
            Proof: Proposition 7.3.5.
        \langle 3 \rangle 2. t = \sec t
            \langle 4 \rangle 1. t \subseteq \operatorname{seg} t
                \langle 5 \rangle 1. Let: s \in t
                \langle 5 \rangle 2. \ s \in \beta
                   PROOF: \beta is a transitive set.
                \langle 5 \rangle 3. \ s \in \operatorname{seg} t
            \langle 4 \rangle 2. seg t \subseteq t
               PROOF: Immediate from definitions.
        \langle 3 \rangle 3. \ \alpha = t
            Proof: Proposition 7.3.3.
        \langle 3 \rangle 4. \ \alpha \in \beta
    \langle 2 \rangle 4. Case: There exists t \in \alpha such that \operatorname{seg} t \cong \beta
       PROOF: \beta \in \alpha similarly.
```

**Proposition 7.3.7** (Z). Any nonempty set S of ordinal numbers has a least element.

```
Proof:
\langle 1 \rangle 1. Pick \beta \in S
\langle 1 \rangle 2. Case: \beta \cap S = \emptyset
  PROOF: Then \beta is least in S.
\langle 1 \rangle 3. Case: \beta \cap S \neq \emptyset
  PROOF: The least element of \beta \cap S is least in S.
Theorem 7.3.8 (ZFC). The class On is well ordered by \in.
Proof:
\langle 1 \rangle 1. Let: \mathbf{E} = \{ (x, y) \in \mathbf{On}^2 \mid x \in y \}
\langle 1 \rangle 2. E is transitive.
  PROOF: If \alpha \in \beta \in \gamma then \alpha \in \gamma because every ordinal is a transitive set.
\langle 1 \rangle 3. E satisfies trichotomy.
  Proof: Proposition 7.3.6.
\langle 1 \rangle 4. E linearly orders On.
  Proof: Proposition 6.2.3.
\langle 1 \rangle 5. E is well founded.
  Proof: Proposition 9.6.2.
Corollary 7.3.8.1 (Burali-Forti Paradox (ZFC)). The class On is a proper
PROOF: If it were a set, it would be a transitive set well-ordered by \in, and hence
a member of itself, contradicting Proposition 1.5.3.
Proposition 7.3.9 (ZFC). Any transitive set of ordinal numbers is an ordinal
number.
PROOF: It is well-ordered by \in by Proposition 6.3.3 and Theorem 7.3.8. \sqcup
Proposition 7.3.10 (Z). \emptyset is an ordinal number.
PROOF: Vacuously, it is a transitive set well-ordered by \in.
Definition 7.3.11. We define 0 = \emptyset.
Proposition 7.3.12 (ZFC). If A is a set of ordinal numbers then \bigcup A is an
ordinal number.
Proof:
\langle 1 \rangle 1. \bigcup A is a transitive set.
  Proof: Proposition 1.6.3.
\langle 1 \rangle 2. \bigcup A is a set of ordinals.
  Proof: Proposition 7.3.5.
\langle 1 \rangle 3. Q.E.D.
  Proof: Proposition 7.3.9.
```

Corollary 7.3.12.1 (ZFC). The poset On is complete.

PROOF: For any nonempty set A of ordinals,  $\bigcup A$  is its supremum.  $\square$ 

**Proposition 7.3.13** (ZFC). Let  $\alpha$  be an ordinal and  $S \subseteq \alpha$ . Then S is well-ordered by  $\in$  and the ordinal of  $(S, \in)$  is  $\leq \alpha$ .

#### Proof:

- $\langle 1 \rangle 1$ . S is well ordered by  $\in$ .
- $\langle 1 \rangle 2$ . Let:  $\beta$  be the ordinal of  $(S, \in)$
- $\langle 1 \rangle 3$ . Let:  $E: S \approx \beta$  be the unique isomorphism.
- $\langle 1 \rangle 4. \ \forall \gamma \in S.E(\gamma) \leq \gamma$ 
  - $\langle 2 \rangle 1$ . Let:  $\gamma \in S$
  - $\langle 2 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall \delta < \gamma. E(\delta) \leq \delta$
  - $\langle 2 \rangle 3$ .  $E(\gamma)$  is the least element of  $\beta$  that is greater than  $E(\delta)$  for all  $\delta < \gamma$
  - $\langle 2 \rangle 4$ .  $\gamma$  is greater than  $E(\delta)$  for all  $\delta < \gamma$
  - $\langle 2 \rangle 5$ .  $E(\gamma) \leq \gamma$
- $\langle 1 \rangle 5. \ \beta \leq \alpha$ 
  - $\langle 2 \rangle 1. \ \forall \gamma < \beta. \gamma < \alpha$ 
    - $\langle 3 \rangle 1$ . Let:  $\gamma < \beta$
    - $\langle 3 \rangle 2$ . Pick  $\delta \in S$  such that  $E(\delta) = \gamma$
    - $\langle 3 \rangle 3. \ \gamma = E(\delta) \le \delta < \alpha$

**Proposition 7.3.14** (ZFC). Let  $\alpha$  be a set. Then the following are equivalent.

- 1.  $\alpha$  is an ordinal.
- 2.  $\alpha$  is a transitive set and, for all  $x, y \in \alpha$ , either x = y or  $x \in y$  or  $y \in x$ .
- 3.  $\alpha$  is a transitive set of transitive sets.

#### Proof:

 $\langle 1 \rangle 1. \ 1 \Rightarrow 2$ 

PROOF: Immediate from definitions.

- $\langle 1 \rangle 2. \ 2 \Rightarrow 3$ 
  - $\langle 2 \rangle 1$ . Assume:  $\alpha$  is a transitive set and, for all  $x,y \in \alpha$ , either x=y or  $x \in y$  or  $y \in x$
  - $\langle 2 \rangle 2$ . Let:  $z \in \alpha$

Prove: z is transitive.

- $\langle 2 \rangle 3$ . Let:  $x \in y \in z$
- $\langle 2 \rangle 4. \ y \in \alpha$
- $\langle 2 \rangle 5. \ x \in \alpha$
- $\langle 2 \rangle 6$ . Either x = z or  $x \in z$  or  $z \in x$
- $\langle 2 \rangle 7. \ x \neq z$

PROOF: We cannot have  $x \in y \in x$  by the Axiom of Regularity.

 $\langle 2 \rangle 8. \ z \notin x$ 

PROOF: We cannot have  $x \in y \in z \in x$  by the Axiom of Regularity.

 $\langle 1 \rangle 3. \ 3 \Rightarrow 1$ 

- $\langle 2 \rangle 1$ . Let: x be a transitive set of transitive sets.
- $\langle 2 \rangle 2$ . Assume: as  $\in$ -induction hypothesis that, for all  $y \in x$ , if y is a transitive set of transitive sets then y is a transitive set of ordinals.
- $\langle 2 \rangle 3$ . Every element of x is an ordinal.
  - $\langle 3 \rangle 1$ . Let:  $y \in x$
  - $\langle 3 \rangle 2$ . y is transitive.
  - $\langle 3 \rangle 3$ . Every element of y is transitive.

PROOF: Since every element of y is an element of x, because x is transitive.

 $\langle 3 \rangle 4$ . y is an ordinal.

Proof:  $\langle 2 \rangle 2$ 

 $\langle 2 \rangle 4$ . Q.E.D.

Proof: Proposition 7.3.9.

**Lemma 7.3.15** (Z). Let A and B be well-ordered sets. If B is an end extension of A then the ordinal of A is  $\leq$  the ordinal of B.

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $\alpha$  be the ordinal of A and  $\beta$  the ordinal of B.
- $\langle 1 \rangle 2$ . Let:  $E_A : A \cong \alpha$  and  $E_B : B \cong \beta$  be the canonical isomorphisms.
- $\langle 1 \rangle 3. \ \forall a \in A.E_A(a) = E_B(a)$ 
  - $\langle 2 \rangle 1$ . Let:  $a \in A$
  - $\langle 2 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall x < a.E_A(x) = E_B(x)$
  - $\langle 2 \rangle 3$ .  $E_A(a)$  is the least ordinal that is greater than  $E_A(x)$  for all x < a
  - $\langle 2 \rangle 4$ .  $E_B(a)$  is the least ordinal that is greater than  $E_B(x)$  for all x < b
  - $\langle 2 \rangle 5. \{ x \in A \mid x <_A a \} = \{ x \in B \mid x <_B a \}$
  - $\langle 2 \rangle 6. \ E_A(a) = E_B(a)$
- $\langle 1 \rangle 4. \ \alpha \subseteq \beta$
- $\langle 1 \rangle 5. \ \alpha \leq \beta$

**Lemma 7.3.16.** Let C be a set of well ordered sets such that, for any  $A, B \in C$ , we have that one of A and B is an end extension of the other. Let  $W = \bigcup C$  under  $x \leq y$  iff there exists  $A \in W$  such that  $x, y \in A$  and  $x \leq y$ . Then W is a well ordered set whose ordinal is the supremum of the ordinals of the members of C.

- $\langle 1 \rangle 1. \leq \text{is reflexive on } W.$ 
  - $\langle 2 \rangle 1$ . Let:  $x \in W$
  - $\langle 2 \rangle 2$ . Pick  $A \in W$  such that  $x \in A$ .
  - $\langle 2 \rangle 3. \ x \leq x$
- $\langle 1 \rangle 2. \leq \text{is antisymmetric on } W.$ 
  - $\langle 2 \rangle 1$ . Let:  $x, y \in W$
  - $\langle 2 \rangle 2$ . Assume:  $x \leq y$  and  $y \leq x$
  - $\langle 2 \rangle 3.$  Pick  $A \in W$  such that  $x,y \in A$  and  $x \leq_A y,$  and  $B \in W$  such that  $x,y \in B$  and  $y \leq_B x$

- $\langle 2 \rangle 4$ . Assume: w.l.o.g. B is an end extension of A
- $\langle 2 \rangle 5$ .  $x \leq_B y$  and  $y \leq_B x$
- $\langle 2 \rangle 6. \ x = y$
- $\langle 1 \rangle 3. \leq \text{is transitive on } W.$ 
  - $\langle 2 \rangle 1$ . Assume:  $x \leq y \leq z$
  - $\langle 2 \rangle 2$ . PICK  $A, B \in W$  such that  $x \leq_A y$  and  $y \leq_B z$
  - $\langle 2 \rangle 3$ . Case: A is an end extension of B.
    - $\langle 3 \rangle 1$ .  $x \leq_A y$  and  $y \leq_A z$
    - $\langle 3 \rangle 2$ .  $x \leq_A z$
    - $\langle 3 \rangle 3. \ x \leq z$
  - $\langle 2 \rangle 4$ . Case: B is an end extension of A.

PROOF: Similar.

- $\langle 1 \rangle 4$ .  $\leq$  is total on W.
  - $\langle 2 \rangle 1$ . Let:  $x, y \in W$
  - $\langle 2 \rangle 2$ . Pick  $A, B \in \mathcal{C}$  such that  $x \in A$  and  $y \in B$
  - $\langle 2 \rangle 3$ . Assume: w.l.o.g. B is an end extension of A
  - $\langle 2 \rangle 4$ .  $x \leq_B y$  or  $y \leq_B x$
  - $\langle 2 \rangle 5$ .  $x \leq_W y$  or  $y \leq_W x$
- $\langle 1 \rangle 5$ . Every nonempty subset of W has a least element.
  - $\langle 2 \rangle 1$ . Let: S be a nonempty subset of W
  - $\langle 2 \rangle 2$ . Pick  $s \in S$
  - $\langle 2 \rangle 3$ . Pick  $A \in \mathcal{C}$  such that  $s \in A$
  - $\langle 2 \rangle 4$ . Let: a be the  $\leq_A$ -least element of  $S \cap A$ Prove: a is least in S
  - $\langle 2 \rangle 5$ . Let:  $x \in S$

Prove:  $a \le x$ 

- $\langle 2 \rangle 6$ . Pick  $B \in \mathcal{C}$  such that  $x \in B$
- $\langle 2 \rangle$ 7. Case: A is an end extension of B
  - $\langle 3 \rangle 1. \ a \leq_A x$
  - $\langle 3 \rangle 2. \ a \leq x$
- $\langle 2 \rangle 8$ . Case: B is an end extension of A
  - $\langle 3 \rangle 1$ . Case:  $x \in A$ 
    - $\langle 4 \rangle 1. \ a \leq_A x$
    - $\langle 4 \rangle 2. \ a \leq x$
  - $\langle 3 \rangle 2$ . Case:  $x \in B A$ 
    - $\langle 4 \rangle 1. \ a \leq_B x$
    - $\langle 4 \rangle 2. \ a \leq x$
- $\langle 1 \rangle 6$ . For all  $A \in \mathcal{C}$ , W is an end extension of A.
  - $\langle 2 \rangle 1$ . For all  $x, y \in A$ , we have  $x \leq_A y$  if and only if  $x \leq_W y$ 
    - $\langle 3 \rangle 1$ . Let:  $x, y \in A$
    - $\langle 3 \rangle 2$ . If  $x \leq_A y$  then  $x \leq_W y$

PROOF: Immediate from definitions.

- $\langle 3 \rangle 3$ . If  $x \leq_W y$  then  $x \leq_A y$ 
  - $\langle 4 \rangle 1$ . Assume:  $x \leq_W y$
  - $\langle 4 \rangle 2$ . PICK  $B \in \mathcal{C}$  such that  $x \leq_B y$
  - $\langle 4 \rangle 3$ . Case: A is an end extension of B

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PROOF: Then  $x \leq_A y$ .

 $\langle 4 \rangle 4$ . Case: B is an end extension of A

PROOF: Then  $x \leq_A y$ .

- $\langle 2 \rangle 2$ . For all  $x \in A$  and  $y \in W A$  we have x < y
  - $\langle 3 \rangle 1$ . Let:  $x \in A$  and  $y \in W A$
  - $\langle 3 \rangle 2$ . PICK  $B \in \mathcal{C}$  such that  $y \in B$
  - $\langle 3 \rangle 3$ . B is an end extension of A
  - $\langle 3 \rangle 4$ .  $x <_B y$
  - $\langle 3 \rangle 5. \ x <_W y$
- $\langle 1 \rangle 7$ . For all  $A \in \mathcal{C}$ , the ordinal of A is  $\leq$  the ordinal of W.

Proof: Lemma 7.3.15.

- $\langle 1 \rangle 8$ . For any ordinal  $\alpha$ , if for all  $A \in \mathcal{C}$  the ordinal of A is  $\leq \alpha$ , then the ordinal of W is  $\leq \alpha$ .
  - $\langle 2 \rangle$ 1. Let:  $\alpha$  be an ordinal.
  - $\langle 2 \rangle 2$ . Assume: for all  $A \in \mathcal{C}$ , the ordinal of A is  $\leq \alpha$
  - $\langle 2 \rangle 3$ . Let:  $\beta$  be the ordinal of W
  - $\langle 2 \rangle 4$ . Let:  $E: W \approx \beta$  be the canonical isomorphism.
  - $\langle 2 \rangle$ 5. Assume: for a contradiction  $\alpha < \beta$
  - $\langle 2 \rangle 6$ . Let:  $a \in W$  be the element with  $E(a) = \alpha$
  - $\langle 2 \rangle$ 7. Pick  $A \in \mathcal{C}$  such that  $a \in A$
  - $\langle 2 \rangle$ 8. Let:  $\gamma$  be the ordinal of A and  $E_A: A \cong \gamma$  be the canonical isomorphism.
  - $\langle 2 \rangle 9$ . For all  $x \in A$  we have  $E_A(x) = E(x)$

PROOF: Transfinite induction on x.

- $\langle 2 \rangle 10. \ E_A(a) = \alpha$
- $\langle 2 \rangle 11. \ \alpha < \gamma$
- $\langle 2 \rangle 12$ . Q.E.D.

PROOF: This contradicts  $\langle 2 \rangle 2$ .

## 7.4 Successors

**Definition 7.4.1** (Successor). The *successor* of a set a is the set  $a^+ := a \cup \{a\}$ .

Proposition 7.4.2 (Z). A set a is a transitive set if and only if

$$\bigcup (a^+) = a .$$

- $\langle 1 \rangle 1$ . If a is a transitive set then  $\bigcup (a^+) = a$ .
  - $\langle 2 \rangle 1$ . Assume: a is a transitive set.
  - $\langle 2 \rangle 2. \ \bigcup (a^+) \subseteq a$ 
    - $\langle 3 \rangle$ 1. Let:  $x \in \bigcup (a^+)$ Prove:  $x \in a$
    - $\langle 3 \rangle 2$ . PICK  $y \in a^+$  such that  $x \in y$ .
    - $\langle 3 \rangle 3. \ y \in a \text{ or } y = a.$

```
\langle 3 \rangle 4. \text{ Case: } y \in a Proof: Then x \in a because a is a transitive set. \langle 3 \rangle 5. \text{ Case: } y = a Proof: Then x \in a immediately. \langle 2 \rangle 3. \ a \subseteq \bigcup (a^+) Proof: Since a \in a^+. \langle 1 \rangle 2. \text{ If } \bigcup (a^+) = a \text{ then } a \text{ is a transitive set.} \langle 2 \rangle 1. \text{ Assume: } \bigcup (a^+) = a \langle 2 \rangle 2. \bigcup a \subseteq a Proof: \bigcup a \subseteq \bigcup (a^+) (Proposition 1.5.9) = a \langle 2 \rangle 3. \ a \text{ is a transitive set.} Proof: Proposition 1.6.2.
```

Proposition 7.4.3. For any set a, we have a is a transitive set if and only if

Proof.

 $\langle 1 \rangle 1$ . If a is a transitive set then  $a^+$  is a transitive set.

PROOF: If a is a transitive set then  $\bigcup (a^+) = a \subseteq a^+$  by Proposition 7.4.2 and so  $a^+$  is a transitive set.

- $\langle 1 \rangle 2$ . If  $a^+$  is a transitive set then a is a transitive set.
  - $\langle 2 \rangle 1$ . Assume:  $a^+$  is a transitive set.
  - $\langle 2 \rangle 2$ . Let:  $x \in y \in a$
  - $\langle 2 \rangle 3. \ x \in y \in a^+$

 $a^+$  is a transitive set.

 $\langle 2 \rangle 4. \ x \in a^+$ 

Proof:  $\langle 2 \rangle 1$ 

 $\langle 2 \rangle 5. \ x \neq a$ 

PROOF: From  $\langle 2 \rangle 2$  and the Axiom of Regularity.

 $\langle 2 \rangle 6. \ x \in a$ 

**Definition 7.4.4.** We write 0 for  $\emptyset$ , 1 for  $\emptyset^+$ , 2 for  $\emptyset^{++}$ , etc.

**Proposition 7.4.5.** For any set A we have  $\mathcal{P}A \approx 2^A$ .

PROOF: The function  $H: \mathcal{P}A \to 2^A$  defined by  $H(S)(a) = \{\emptyset\}$  if  $a \in S$  and  $\emptyset$  if  $a \notin S$  is a bijection.  $\square$ 

**Proposition 7.4.6.** For any ordinal number  $\alpha$  we have  $\alpha^+$  is an ordinal number.

Proof:

 $\langle 1 \rangle 1$ .  $\alpha^+$  is a transitive set.

Proof: Proposition 7.4.3.

 $\langle 1 \rangle 2$ .  $\alpha^+$  is well-ordered by  $\in$ .

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 $\langle 2 \rangle 1$ . For all  $x, y, z \in \alpha^+$ , if  $x \in y \in z$  then  $x \in z$ 

 $\langle 3 \rangle 1$ . Case:  $z = \alpha$ 

PROOF: Then  $x \in \alpha$  since  $\alpha$  is a transitive set.

 $\langle 3 \rangle 2$ . Case:  $z \in \alpha$ 

PROOF: Then  $x \in z$  since  $\alpha$  is well-ordered by  $\in$ .

 $\langle 2 \rangle 2$ . For all  $x, y \in \alpha^+$  we have  $x \in y$  or x = y or  $y \in x$ 

 $\langle 3 \rangle 1$ . Case:  $x, y \in \alpha$ 

PROOF: The result follows because  $\alpha$  is well-ordered by  $\in$ .

 $\langle 3 \rangle 2$ . Case:  $x \in \alpha, y = \alpha$ 

PROOF: Then  $x \in y$ .

 $\langle 3 \rangle 3$ . Case:  $x = \alpha, y \in \alpha$ 

PROOF: Then  $y \in x$ .

 $\langle 3 \rangle 4$ . Case:  $x = \alpha, y = \alpha$ 

PROOF: Then x = y.

 $\langle 2 \rangle 3$ . Every nonempty subset of  $\alpha^+$  has an  $\in$ -least element.

 $\langle 3 \rangle 1$ . Let:  $S \subseteq \alpha^+$  be nonempty

 $\langle 3 \rangle 2$ . Case:  $S = \{\alpha\}$ 

PROOF:  $\alpha$  is least in S.

 $\langle 3 \rangle 3$ . Case:  $S \neq \{\alpha\}$ 

 $\langle 4 \rangle 1$ .  $S - \{\alpha\}$  is a nonempty subset of  $\alpha$ 

 $\langle 4 \rangle 2$ . Let:  $\beta$  be least in  $S - \{\alpha\}$ 

 $\langle 4 \rangle 3$ .  $\beta$  is least in S.

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**Proposition 7.4.7.** For ordinals  $\alpha$  and  $\beta$ , if  $\alpha^+ = \beta^+$  then  $\alpha = \beta$ .

PROOF: If  $\alpha^+ = \beta^+$  then  $\alpha = \bigcup (\alpha^+)$ (Proposition 7.4.2)  $=\bigcup(\beta^+)$  $=\beta$ (Proposition 7.4.2)

**Proposition 7.4.8.** For ordinals  $\alpha$  and  $\beta$ , we have  $\alpha < \beta$  if and only if  $\alpha^+ < \beta$  $\beta^+$ .

Proof:

$$\alpha < \beta \Leftrightarrow \alpha^+ \le \beta$$
$$\Leftrightarrow \alpha^+ < \beta^+$$

**Definition 7.4.9** (Successor Ordinal). An ordinal  $\alpha$  is a successor ordinal iff  $\alpha = \beta^+$  for some  $\beta$ .

**Definition 7.4.10** (Limit Ordinal). A *limit ordinal* is an ordinal that is neither 0 nor a successor ordinal.

**Proposition 7.4.11.** *If*  $\lambda$  *is a limit ordinal and*  $\beta < \lambda$  *then*  $\beta^+ < \lambda$ .

PROOF: Since  $\beta^+ \leq \lambda$  and  $\beta^+ \neq \lambda$ .  $\square$ 

## 7.5 The Well-Ordering Theorem and Zorn's Lemma

**Theorem 7.5.1** (Hartogs). For any set A, there exists an ordinal not dominated by A.

#### Proof:

- $\langle 1 \rangle$ 1. Let:  $\alpha$  be the class of all ordinals  $\beta$  such that  $\beta \preccurlyeq A$  Prove:  $\alpha$  is a set.
- $\langle 1 \rangle 2$ . Let:  $W = \{(B, R) \mid B \subseteq A, R \text{ is a well ordering on } B\}$
- $\langle 1 \rangle 3$ .  $\alpha$  is the class of the ordinals of the elements of W.
  - $\langle 2 \rangle 1$ . For all  $(B,R) \in W$ , the ordinal of (B,R) is in  $\alpha$ .
    - $\langle 3 \rangle 1$ . Let:  $(B, R) \in W$
    - $\langle 3 \rangle 2$ . Let:  $\beta$  be the ordinal of (B, R)
    - $\langle 3 \rangle 3$ . Let:  $E : B \cong \beta$  be the canonical isomorphism.
    - $\langle 3 \rangle 4$ . Let:  $i: B \hookrightarrow A$  be the inclusion
    - $\langle 3 \rangle 5.$   $i \circ E^{-1}$  is an injection  $\beta \to A$
    - $\langle 3 \rangle 6. \ \beta \in \alpha$
  - $\langle 2 \rangle 2$ . For all  $\beta \in \alpha$ , there exists  $(B,R) \in W$  such that  $\beta$  is the ordinal number of (B,R).
    - $\langle 3 \rangle 1$ . Let:  $\beta \in \alpha$
    - $\langle 3 \rangle 2$ . Pick an injection  $f: \beta \to A$
    - $\langle 3 \rangle 3$ . Define  $\leq$  on ran f by  $f(x) \leq f(y)$  iff  $x \leq y$
    - $\langle 3 \rangle 4$ .  $(\operatorname{ran} f, \leq) \in W$
    - $\langle 3 \rangle 5$ .  $\beta$  is the ordinal number of (ran f, <)
- $\langle 1 \rangle 4$ .  $\alpha$  is a set.

PROOF: By an Axiom of Replacement.

 $\langle 1 \rangle 5$ .  $\alpha$  is an ordinal.

PROOF: It is a transitive set of ordinals.

 $\langle 1 \rangle 6. \ \alpha \not \leq A$ 

PROOF: Since  $\alpha \notin \alpha$ .

**Theorem 7.5.2** (Numeration Theorem). Every set is equinumerous with some ordinal.

#### PROOF:

- $\langle 1 \rangle 1$ . Let: A be any set.
- $\langle 1 \rangle 2$ . PICK an ordinal  $\alpha$  not dominated by A.
- $\langle 1 \rangle 3$ . Pick a choice function G for A.
- $\langle 1 \rangle 4$ . Pick  $e \notin A$
- $\langle 1 \rangle 5$ . Let:  $F: \alpha \to A \cup \{e\}$  by transfinite recursion:

$$F(\gamma) = \begin{cases} G(A - F(\{\delta \mid \delta < \gamma\}) & \text{if } A - F(\{\delta \mid \delta < \gamma\}) \neq \emptyset \\ e & \text{if } A - F(\{\delta \mid \delta < \gamma\}) = \emptyset \end{cases}$$

- $\langle 1 \rangle 6. \ e \in \operatorname{ran} F$ 
  - $\langle 2 \rangle 1$ . Assume: for a contradiction  $e \notin \operatorname{ran} F$
  - $\langle 2 \rangle 2$ . F is an injection  $\alpha \to A$ .
    - $\langle 3 \rangle 1$ . Let:  $\beta, \gamma \in \alpha$  with  $\beta \neq \gamma$

```
Prove: F(\beta) \neq F(\gamma)
        \langle 3 \rangle 2. Assume: w.l.o.g. \beta < \gamma
        \langle 3 \rangle 3. \ F(\gamma) \in A - F(\{\delta \mid \delta < \gamma\})
        \langle 3 \rangle 4. \ F(\gamma) \notin F(\{\delta \mid \delta < \gamma\})
        \langle 3 \rangle 5. \ F(\gamma) \neq F(\beta)
    \langle 2 \rangle 3. Q.E.D.
        PROOF: This contradicts \langle 1 \rangle 2.
\langle 1 \rangle 7. Let: \delta be least such that F(\delta) = e
\langle 1 \rangle 8. \ F \upharpoonright \delta : \delta \approx A
Theorem 7.5.3 (Well-Ordering Theorem). Any set can be well ordered.
Proof:
\langle 1 \rangle 1. Pick an ordinal \delta and a bijection F: A \approx \delta
\langle 1 \rangle 2. Define \leq on A by F(x) \leq F(y) iff x \leq y for x, y \in \delta
\langle 1 \rangle 3. \leq \text{is a well ordering on } A.
Theorem 7.5.4 (Zorn's Lemma). Let A be a set such that, for every chain
\mathcal{B} \subseteq \mathcal{A}, we have \bigcup \mathcal{B} \in \mathcal{A}. Then \mathcal{A} has a maximal element.
Proof:
\langle 1 \rangle 1. PICK a well ordering \langle 0 \rangle on \mathcal{A}.
\langle 1 \rangle 2. Let: F: \mathcal{A} \to 2 be the function defined by transfinite recursion by:
                        F(A) = \begin{cases} 1 & \text{if } A \text{ includes every set } B < A \text{ for which } F(B) = 1 \\ 0 & \text{otherwise} \end{cases}
\langle 1 \rangle 3. Let: \mathcal{C} = \{ A \in \mathcal{A} \mid F(A) = 1 \}
           PROVE: \bigcup \mathcal{C} is a maximal element of \mathcal{A}
\langle 1 \rangle 4. For all A \in \mathcal{A}, we have A \in \mathcal{C} iff \forall B < A.B \in \mathcal{C} \Rightarrow B \subseteq A
\langle 1 \rangle 5. C is a chain.
    \langle 2 \rangle 1. Let: A, A' \in \mathcal{C}
    \langle 2 \rangle 2. Assume: w.l.o.g. A \leq A'
    \langle 2 \rangle 3. \ A \subseteq A'
        Proof: By \langle 1 \rangle 4
\langle 1 \rangle 6. \bigcup \mathcal{C} \in \mathcal{A}
\langle 1 \rangle 7. \bigcup C is maximal in A.
    \langle 2 \rangle 1. Let: A \in \mathcal{A} and \bigcup \mathcal{C} \subseteq A
    \langle 2 \rangle 2. A \in \mathcal{C}
        PROOF: By \langle 1 \rangle 4 since \forall B \in \mathcal{C}.B \subseteq A.
    \langle 2 \rangle 3. \ A \subseteq \bigcup \mathcal{C}
     \langle 2 \rangle 4. A = \bigcup \mathcal{C}
```

**Proposition 7.5.5** (Teichmüller-Tukey Lemma). Let A be a nonempty set such that, for every B, we have  $B \in A$  if and only if every finite subset of B is a member of A. Then A has a maximal element.

Proof:

```
\langle 1 \rangle 1. For every chain \mathcal{B} \subseteq \mathcal{A}, we have \bigcup \mathcal{B} \in \mathcal{A} \langle 2 \rangle 1. Let: \mathcal{B} \subseteq \mathcal{A} be a chain. \langle 2 \rangle 2. Every finite subset of \bigcup \mathcal{B} is a member of \mathcal{A}. \langle 3 \rangle 1. Let: C be a finite subset of \bigcup \mathcal{B}. \langle 3 \rangle 2. Pick \mathcal{B} \in \mathcal{B} such that C \subseteq \mathcal{B}. \langle 3 \rangle 3. \mathcal{B} \in \mathcal{A} \langle 3 \rangle 4. Every finite subset of \mathcal{B} is in \mathcal{A}. \langle 3 \rangle 5. C \in \mathcal{A} \langle 2 \rangle 3. \bigcup \mathcal{B} \in \mathcal{A}. \langle 1 \rangle 2. Q.E.D. Proof: Zorn's lemma.
```

**Theorem Schema 7.5.6.** For any class **A**, there exists a class **F** such that the following is a theorem:

If A is a proper class of ordinals, then  $F : \mathbf{On} \to A$  is an order isomorphism.

#### **PROOF**

- $\langle 1 \rangle 1$ . Define  $\mathbf{F} : \mathbf{On} \to \mathbf{A}$  by transfinite recursion as follows:  $\mathbf{F}(\alpha)$  is the least element of  $\mathbf{A}$  that is different from  $\mathbf{F}(\beta)$  for all  $\beta < \alpha$ .
- $\langle 1 \rangle 2$ . For all  $\alpha, \beta \in \mathbf{On}$ , if  $\alpha < \beta$  then  $\mathbf{F}(\alpha) < \mathbf{F}(\beta)$ PROOF: We have  $\mathbf{F}(\alpha) \neq \mathbf{F}(\beta)$  by the definition of  $\mathbf{F}(\beta)$ , and  $\mathbf{F}(\beta) \not < \mathbf{F}(\alpha)$  by the leastness of  $\mathbf{F}(\alpha)$ .
- $\langle 1 \rangle 3$ . **F** is surjective.
  - $\langle 2 \rangle 1$ . Let:  $\alpha \in \mathbf{A}$
  - $\langle 2 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall \beta \in \mathbf{A}$ , if  $\beta < \alpha$  then there exists  $\gamma$  such that  $\beta = \mathbf{F}(\gamma)$ .
  - $\langle 2 \rangle 3$ . Let:  $\gamma = \{ \delta \in \mathbf{On} \mid \mathbf{F}(\delta) < \alpha \}$
  - $\langle 2 \rangle 4$ .  $\gamma$  is a set.

PROOF: Axiom of Replacement applied to  $\alpha$ .

 $\langle 2 \rangle$ 5.  $\gamma$  is a transitive set.

PROOF: If  $\mathbf{F}(\delta) < \alpha$  and  $\epsilon < \delta$  then  $\mathbf{F}(\epsilon) < \alpha$  by  $\langle 1 \rangle 2$ .

 $\langle 2 \rangle 6$ .  $\gamma$  is an ordinal.

Proof: Proposition 7.3.9.

- $\langle 2 \rangle 7$ .  $\mathbf{F}(\gamma) = \alpha$ 
  - $\langle 3 \rangle 1$ .  $\mathbf{F}(\gamma)$  is the least element of **A** different from  $\mathbf{F}(\delta)$  for all  $\delta < \gamma$
  - $\langle 3 \rangle 2$ .  $\mathbf{F}(\gamma)$  is the least element of  $\mathbf{A}$  different from x for all  $x \in \mathbf{A}$  with  $x < \alpha$
  - $\langle 3 \rangle 3. \ \mathbf{F}(\gamma) = \alpha$

## 7.6 Ordinal Operations

**Definition 7.6.1** (Ordinal Operation). An *ordinal operation* is a function  $\mathbf{On} \to \mathbf{On}$ .

**Definition 7.6.2** (Continuous). An ordinal operation  $\mathbf{T}: \mathbf{On} \to \mathbf{On}$  is *continuous* iff, for every limit ordinal  $\lambda$ , we have  $\mathbf{T}(\lambda) = \bigcup_{\alpha < \lambda} \mathbf{T}(\alpha)$ .

**Definition 7.6.3** (Normal). An ordinal operation is *normal* iff it is continuous and strictly monotone.

**Proposition Schema 7.6.4.** For any class **T**, the following is a theorem.

If **T** is a continuous ordinal operation and  $\forall \gamma. \mathbf{T}(\gamma) < \mathbf{T}(\gamma^+)$ , then **T** is normal.

Proof:

 $\langle 1 \rangle 1$ . Let:  $P[\beta]$  be the property  $\forall \gamma < \beta . \mathbf{T}(\gamma) < \mathbf{T}(\beta)$ 

 $\langle 1 \rangle 2$ . P[0]

Proof: Vacuous.

 $\langle 1 \rangle 3$ . For any ordinal  $\gamma$ , if  $P[\gamma]$  then  $P[\gamma^+]$ 

 $\langle 2 \rangle 1$ . Assume:  $P[\gamma]$ 

 $\langle 2 \rangle 2$ . Let:  $\delta < \gamma^+$ 

 $\langle 2 \rangle 3$ . Case:  $\delta < \gamma$ 

PROOF: Then  $\mathbf{T}(\delta) < \mathbf{T}(\gamma) < \mathbf{T}(\gamma^+)$ .

 $\langle 2 \rangle 4$ . Case:  $\delta = \gamma$ 

PROOF: Then  $\mathbf{T}(\delta) = \mathbf{T}(\gamma) < \mathbf{T}(\gamma^+)$ .

 $\langle 1 \rangle 4$ . For any limit ordinal  $\lambda$ , if  $\forall \gamma < \lambda . P[\gamma]$  then  $P[\lambda]$ .

 $\langle 2 \rangle 1$ . Assume:  $\forall \gamma < \lambda . P[\gamma]$ 

 $\langle 2 \rangle 2$ . Let:  $\delta < \lambda$ 

 $\langle 2 \rangle 3$ .  $\mathbf{T}(\delta) < \mathbf{T}(\lambda)$ 

Proof:

$$\mathbf{T}(\delta) < \mathbf{T}(\delta^{+})$$

$$\leq \bigcup_{\epsilon < \lambda} \mathbf{T}(\epsilon)$$

$$= \mathbf{T}(\lambda)$$

**Proposition Schema 7.6.5.** For any class **T**, the following is a theorem:

Assume **T** is a normal ordinal operation. For every ordinal  $\alpha$ , we have  $\alpha \leq \mathbf{T}(\alpha)$ .

Proof:

 $\langle 1 \rangle 1$ . Let:  $\gamma$  be an ordinal.

 $\langle 1 \rangle 2$ . Assume: as induction hypothesis  $\forall \delta < \gamma. \mathbf{T}(\delta) \geq \delta$ 

 $\langle 1 \rangle 3$ . For all  $\delta < \gamma$  we have  $\delta < \mathbf{T}(\gamma)$ 

PROOF: **T** is strictly monotone.

 $\langle 1 \rangle 4. \ \gamma \leq \mathbf{T}(\gamma)$ 

**Proposition Schema 7.6.6.** For any class T, the following is a theorem:

Assume **T** is a normal ordinal operation. For any ordinal  $\beta \geq \mathbf{T}(0)$ , there exists a greatest ordinal  $\gamma$  such that  $\mathbf{T}(\gamma) \leq \beta$ .

```
Proof:
```

- $\langle 1 \rangle 1$ . There exists  $\gamma$  such that  $\mathbf{T}(\gamma) > \beta$ 
  - $\langle 2 \rangle 1$ . For all  $\gamma$  we have  $\mathbf{T}(\gamma) \geq \gamma$

Proof: Proposition 7.6.5.

- $\langle 2 \rangle 2$ .  $\mathbf{T}(\beta^+) > \beta$
- $\langle 1 \rangle 2$ . Let:  $\delta$  be least such that  $\mathbf{T}(\delta) > \beta$
- $\langle 1 \rangle 3$ .  $\delta$  is a successor ordinal.
  - $\langle 2 \rangle 1. \ \delta \neq 0$

PROOF: Since  $\mathbf{T}(0) \leq \beta$ .

- $\langle 2 \rangle 2$ .  $\delta$  is not a limit ordinal.
  - $\langle 3 \rangle 1$ . Assume: for a contradiction  $\delta$  is a limit ordinal.
  - $\langle 3 \rangle 2. \ \beta < \bigcup_{\epsilon < \delta} \mathbf{T}(\epsilon)$

PROOF: T is continuous.

- $\langle 3 \rangle 3$ . There exists  $\epsilon < \delta$  such that  $\beta < \mathbf{T}(\epsilon)$
- $\langle 3 \rangle 4$ . Q.E.D.

PROOF: This contradicts the minimality of  $\delta$ .

- $\langle 1 \rangle 4$ . Let:  $\delta = \gamma^+$
- $\langle 1 \rangle 5$ .  $\gamma$  is greatest such that  $\mathbf{T}(\gamma) \leq \beta$

**Theorem Schema 7.6.7.** For any class **T**, the following is a theorem:

Assume that T is a normal ordinal operation. For any nonempty set of ordinals S, we have

$$\mathbf{T}(\sup S) = \sup_{\alpha \in S} \mathbf{T}(\alpha) .$$

## Proof:

 $\langle 1 \rangle 1. \ \forall \alpha \in S. \mathbf{T}(\alpha) \leq \mathbf{T}(\sup S)$ 

PROOF: Since T is monotone.

- $\langle 1 \rangle 2$ . For any ordinal  $\beta$ , if  $\forall \alpha \in S. \mathbf{T}(\alpha) \leq \beta$ , then  $\mathbf{T}(\sup S) \leq \beta$ 
  - $\langle 2 \rangle 1$ . Let:  $\beta$  be an ordinal.
  - $\langle 2 \rangle 2$ . Let:  $\gamma = \sup S$
  - $\langle 2 \rangle 3$ . Assume:  $\forall \alpha \in S.\mathbf{T}(\alpha) \leq \beta$
  - $\langle 2 \rangle 4$ . Case:  $\gamma$  is 0 or a successor ordinal

PROOF: Then we must have  $\gamma \in S$  so  $\mathbf{T}(\gamma) \leq \beta$  from  $\langle 2 \rangle 3$ .

- $\langle 2 \rangle$ 5. Case:  $\gamma$  is a limit ordinal
  - $\langle 3 \rangle 1$ .  $\mathbf{T}(\gamma) = \sup_{\alpha < \gamma} \mathbf{T}(\alpha)$

PROOF: T is continuous.

- $\langle 3 \rangle 2$ . Assume: for a contradiction  $\beta < \mathbf{T}(\gamma)$
- $\langle 3 \rangle 3$ . Pick  $\alpha < \gamma$  such that  $\beta < \mathbf{T}(\alpha)$

Proof:  $\langle 3 \rangle 1$ ,  $\langle 3 \rangle 2$ 

 $\langle 3 \rangle 4$ . PICK  $\alpha' \in S$  such that  $\alpha < \alpha'$ 

Proof:  $\langle 2 \rangle 2$ ,  $\langle 3 \rangle 3$ 

 $\langle 3 \rangle 5. \ \beta < \mathbf{T}(\alpha') \leq \beta$ 

PROOF: **T** is strictly monotone,  $\langle 3 \rangle 3$ ,  $\langle 3 \rangle 4$ ,  $\langle 2 \rangle 3$ .

 $\langle 3 \rangle 6$ . Q.E.D.

Proof: This is a contradiction.

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**Proposition 7.6.8.** For any classes **A** and **T**, the following is a theorem:

Assume **A** is a proper class of ordinals such that, for every set  $S \subseteq \mathbf{A}$ , we have  $\bigcup S \in \mathbf{A}$ . Assume **T** is the unique order isomorphism  $\mathbf{On} \cong \mathbf{A}$ . Then **T** is normal.

#### Proof:

 $\langle 1 \rangle 1$ . **T** is strictly monotone.

PROOF: Since it is an order isomorphism.

- $\langle 1 \rangle 2$ . **T** is continuous.
  - $\langle 2 \rangle 1$ . Let:  $\lambda$  be a limit ordinal.
  - $\langle 2 \rangle 2$ .  $\mathbf{T}'(\lambda)$  is the least member of **A** that is greater than  $\mathbf{T}'(\alpha)$  for all  $\alpha < \lambda$
- $\langle 2 \rangle 3. \ \mathbf{T}'(\lambda) = \sup_{\alpha < \lambda} \mathbf{T}'(\alpha)$

**Proposition Schema 7.6.9.** For any class **T**, the following is a theorem:

If **T** is a normal ordinal operation, then for any limit ordinal  $\lambda$ , we have  $\mathbf{T}(\lambda)$  is a limit ordinal.

### Proof:

```
\langle 1 \rangle 1. \mathbf{T}(\lambda) \neq 0
```

PROOF: Since  $0 \leq \mathbf{T}(0) < \mathbf{T}(\lambda)$ .

- $\langle 1 \rangle 2$ .  $\mathbf{T}(\lambda)$  is not a successor ordinal.
- $\langle 2 \rangle 1$ . Assume: for a contradiction  $\mathbf{T}(\lambda) = \alpha^+$
- $\langle 2 \rangle 2$ .  $\alpha < \mathbf{T}(\lambda) = \sup_{\beta < \lambda} \mathbf{T}(\beta)$
- $\langle 2 \rangle 3$ . PICK  $\beta < \lambda$  such that  $\alpha < \mathbf{T}(\beta)$
- $\langle 2 \rangle 4. \ \alpha^+ \leq \mathbf{T}(\beta) < \mathbf{T}(\lambda)$
- $\langle 2 \rangle$ 5. Q.E.D.

PROOF: This is a contradiction.

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## 7.7 Ordinal Arithmetic

## 7.7.1 Addition

**Definition 7.7.1.** Let A and B be disjoint well-ordered sets. The *concatenation* of A and B is the set  $A \cup B$  under the relation:

- if  $a, a' \in A$  then  $a \leq a'$  iff  $a \leq a'$  in A
- if  $b, b' \in B$  then  $b \leq b'$  iff  $b \leq b'$  in B
- if  $a \in A$  and  $b \in B$  then  $a \le b$  and  $b \not\le a$ .

**Proposition 7.7.2.** If A and B are disjoint well-ordered sets, then their concatenation is well-ordered.

Proof:

 $\langle 1 \rangle 1. \leq \text{is reflexive.}$ 

PROOF: For all  $a \in A$  we have  $a \le a$ , and for all  $b \in B$  we have  $b \le b$ .

- $\langle 1 \rangle 2$ .  $\leq$  is antisymmetric.
  - $\langle 2 \rangle 1$ . Assume:  $x \leq y \leq x$
  - $\langle 2 \rangle 2$ . Case:  $x, y \in A$

PROOF: Then x = y since the order on A is antisymmetric.

 $\langle 2 \rangle 3$ . Case:  $x \in A$  and  $y \in B$ 

PROOF: This is impossible as it would imply  $y \not\leq x$ .

 $\langle 2 \rangle 4$ . Case:  $x \in B$  and  $y \in A$ 

PROOF: This is impossible as it would imply  $x \not\leq y$ .

 $\langle 2 \rangle$ 5. Case:  $x, y \in B$ 

PROOF: Then x = y since the order on B is antisymmetric.

- $\langle 1 \rangle 3. \leq \text{is transitive.}$ 
  - $\langle 2 \rangle 1$ . Assume:  $x \leq y \leq z$
  - $\langle 2 \rangle 2$ . Case:  $x, z \in A$

PROOF: In this case  $y \in A$  since  $y \le z$ , and so  $x \le z$  since the order on A is transitive.

 $\langle 2 \rangle 3$ . Case:  $x \in A$  and  $z \in B$ 

PROOF: Then  $x \leq z$  immediately.

 $\langle 2 \rangle 4$ . Case:  $x \in B$  and  $z \in A$ 

PROOF: This is impossible because we have  $y \notin A$  since  $x \leq y$  and  $y \notin B$  since  $y \leq z$ .

 $\langle 2 \rangle$ 5. Case:  $x, z \in B$ 

PROOF: In this case  $y \in B$  since  $x \le y$ , and so  $x \le z$  since the order on B is transitive.

- $\langle 1 \rangle 4. \leq \text{is total.}$ 
  - $\langle 2 \rangle 1$ . Let:  $x, y \in A \cup B$
  - $\langle 2 \rangle 2$ . Case:  $x, y \in A$

PROOF: Then  $x \leq y$  or  $y \leq x$  because the order on A is total.

 $\langle 2 \rangle 3$ . Case:  $x \in A$  and  $y \in B$ 

PROOF: Then  $x \leq y$ .

 $\langle 2 \rangle 4$ . Case:  $x \in B$  and  $y \in A$ 

PROOF: Then  $y \leq x$ .

 $\langle 2 \rangle$ 5. Case:  $x, y \in B$ 

PROOF: Then  $x \leq y$  or  $y \leq x$  because the order on B is total.

- $\langle 1 \rangle 5$ . Every nonempty subset of  $A \cup B$  has a least element.
  - $\langle 2 \rangle 1$ . Let: S be a nonempty subset of  $A \cup B$
  - $\langle 2 \rangle 2$ . Case:  $S \cap A = \emptyset$

PROOF: Then  $S \subseteq B$  and so S has a least element.

 $\langle 2 \rangle 3$ . Case:  $S \cap A \neq \emptyset$ 

PROOF: The least element of  $S \cap A$  is the least element of S.

**Definition 7.7.3** (Ordinal Addition). Let  $\alpha$  and  $\beta$  be ordinal numbers. Then  $\alpha + \beta$  is the ordinal number of the concatenation of A and B, where A is any well ordered set with ordinal  $\alpha$  and B is any well ordered set with ordinal  $\beta$ .

**Theorem 7.7.4** (Associative Law for Addition). For any ordinals  $\rho$ ,  $\sigma$  and  $\tau$ , we have

$$\rho + (\sigma + \tau) = (\rho + \sigma) + \tau .$$

PROOF: Given disjoint well ordered sets A, B and C, the concatenation of A with (the concatenation of B and C) is the same as the concatenation of (the concatenation of A and B) and C.  $\square$ 

**Theorem 7.7.5.** For any ordinal  $\rho$  we have

$$\rho + 0 = 0 + \rho = \rho$$
.

PROOF: For any well ordered set A, the concatenation of A with  $\emptyset$  is A, and the concatenation of  $\emptyset$  with A is A.  $\square$ 

**Theorem 7.7.6.** For any ordinal  $\alpha$  we have  $\alpha + 1 = \alpha^+$ .

PROOF: Since  $\alpha^+$  is the concatenation of  $\alpha$  and  $\{\alpha\}$ .  $\square$ 

**Theorem 7.7.7.** For any ordinal  $\alpha$ , the operation that maps  $\beta$  to  $\alpha + \beta$  is normal.

Proof:

- $\langle 1 \rangle 1$ . For any limit ordinal  $\lambda$ , we have  $\alpha + \lambda = \sup_{\beta < \lambda} (\alpha + \beta)$ .
  - $\langle 2 \rangle$ 1. Let:  $\lambda$  be a limit ordinal.
  - $\langle 2 \rangle 2$ .  $(\{0\} \times \alpha) \cup (\{1\} \times \lambda) = \bigcup_{\beta \in \lambda} ((\{0\} \times \alpha) \cup (\{1\} \times \beta))$ , where the order on the right hand side is as in Lemma 7.3.16.

Proof:

$$\begin{split} (\{0\} \times \alpha) \cup (\{1\} \times \lambda) &= (\{0\} \times \alpha) \cup (\{1\} \times \bigcup_{\beta < \lambda} \beta) \\ &= (\{0\} \times \alpha) \cup \bigcup_{\beta < \lambda} (\{1\} \times \beta) \\ &= \bigcup_{\beta < \lambda} ((\{0\} \times \alpha) \cup (\{1\} \times \beta)) \end{split}$$

 $\langle 1 \rangle$ 2. For any ordinal  $\beta$  we have  $\alpha + \beta < \alpha + \beta^+$ PROOF: Since  $\alpha + \beta^+ = \alpha + \beta + 1 = (\alpha + \beta)^+$ .

**Corollary 7.7.7.1.** For any ordinals  $\alpha$ ,  $\beta$ ,  $\gamma$ , we have  $\beta < \gamma$  if and only if  $\alpha + \beta < \alpha + \gamma$ .

**Corollary 7.7.7.2** (Left Cancellation for Addition). For any ordinals  $\alpha$ ,  $\beta$  and  $\gamma$ , if  $\alpha + \beta = \alpha + \gamma$  then  $\beta = \gamma$ .

**Theorem 7.7.8.** For any ordinals  $\alpha$ ,  $\beta$ ,  $\gamma$ , if  $\beta \leq \gamma$  then  $\beta + \alpha \leq \gamma + \alpha$ .

PROOF: Transfinite induction on  $\alpha$ .

**Theorem 7.7.9** (Subtraction Theorem). Let  $\alpha$  and  $\beta$  be ordinals with  $\alpha \leq \beta$ . Then there exists a unique ordinal  $\delta$  such that  $\alpha + \delta = \beta$ .

Proof:

- $\langle 1 \rangle 1$ . For all ordinals  $\alpha$  and  $\beta$  with  $\alpha \leq \beta$ , there exists  $\delta$  such that  $\alpha + \delta = \beta$ 
  - $\langle 2 \rangle 1$ . Let:  $\alpha$  and  $\beta$  be ordinals with  $\alpha \leq \beta$
  - $\langle 2 \rangle 2$ . Let:  $\delta$  be the greatest ordinal such that  $\alpha + \delta \leq \beta$

Proof: Proposition 7.6.6.

 $\langle 2 \rangle 3. \ \alpha + \delta = \beta$ 

PROOF: If  $\alpha + \delta < \beta$  then  $\alpha + \delta + 1 \le \beta$  contradicting the greatestness of  $\delta$ .  $\langle 1 \rangle 2$ . Q.E.D.

PROOF: Uniqueness follows from the Left Cancellation Law.

## 7.7.2 Multiplication

**Definition 7.7.10** (Ordinal Multiplication). Let  $\alpha$  and  $\beta$  be ordinal numbers. Then  $\alpha\beta$  is the ordinal number of  $A \times B$  under the lexicographic order, where A is any well ordered set with ordinal  $\alpha$  and B is any well ordered set with ordinal  $\beta$ .

This is well defined by Proposition 6.3.5.

**Theorem 7.7.11** (Associative Law). For any ordinals  $\rho$ ,  $\sigma$  and  $\tau$ , we have

$$\rho(\sigma\tau) = (\rho\sigma)\tau \ .$$

PROOF: Let A, B and C be well ordered sets with ordinals  $\rho$ ,  $\sigma$  and  $\tau$ . Then both  $\rho(\sigma\tau)$  and  $(\rho\sigma)\tau$  are the ordinal of  $A\times B\times C$  under  $(a,b,c)\leq (a',b',c')\Leftrightarrow a\leq a'\vee (a=a'\wedge b\leq b')\vee (a=a'\wedge b=b'\wedge c\leq c')$ .

**Theorem 7.7.12** (Left Distributive Law). For any ordinals  $\rho$ ,  $\sigma$  and  $\tau$ , we have

$$\rho(\sigma + \tau) = \rho\sigma + \rho\tau$$

PROOF: Let A, B and C be well ordered sets with ordinals  $\rho$ ,  $\sigma$  and  $\tau$  and with  $B \cap C = \emptyset$ . Then both  $\rho(\sigma + \tau)$  and  $\rho\sigma + \rho\tau$  are the ordinal of  $A \times (B \cup C)$  under the lexicographic ordering.  $\square$ 

**Theorem 7.7.13.** For any ordinal  $\rho$  we have  $\rho 0 = 0 \rho = 0$ .

PROOF: For any well ordered set A we have  $A \times \emptyset = \emptyset \times A = \emptyset$ .  $\square$ 

**Theorem 7.7.14.** For any ordinal  $\rho$  we have  $\rho 1 = 1\rho = \rho$ .

Proof: Easy.  $\square$ 

**Theorem 7.7.15.** For any ordinals  $\rho$  and  $\sigma$ , if  $\rho\sigma = 0$  then  $\rho = 0$  or  $\sigma = 0$ .

PROOF: If  $A \times B = \emptyset$  then  $A = \emptyset$  or  $B = \emptyset$ .  $\square$ 

**Theorem 7.7.16.** For any non-zero ordinal  $\alpha$ , the operation that maps  $\beta$  to  $\alpha\beta$  is normal.

Proof:

 $\langle 1 \rangle 1$ . For any limit ordinal  $\lambda$ , we have  $\alpha \lambda = \bigcup_{\beta < \lambda} \alpha \beta$ 

 $\langle 2 \rangle 1$ . Let:  $\lambda$  be a limit ordinal

 $\langle 2 \rangle 2$ .  $\alpha \times \lambda = \bigcup_{\beta < \lambda} (\alpha \times \beta)$  as well-ordered sets

 $\langle 1 \rangle 2$ . For any ordinal  $\beta$  we have  $\alpha \beta < \alpha \beta^+$ 

PROOF:  $\alpha \beta^+ = \alpha \beta + \alpha > \alpha \beta$ 

**Corollary 7.7.16.1.** For any ordinals  $\alpha$ ,  $\beta$ ,  $\gamma$ , if  $\alpha \neq 0$  then  $\beta < \gamma$  if and only if  $\alpha\beta < \alpha\gamma$ .

**Corollary 7.7.16.2** (Left Cancellation for Multiplication). For any ordinals  $\alpha$ ,  $\beta$ ,  $\gamma$ , if  $\alpha \neq 0$  and  $\alpha\beta = \alpha\gamma$  then  $\beta = \gamma$ .

**Theorem 7.7.17.** For any ordinals  $\alpha$ ,  $\beta$  and  $\gamma$ , if  $\beta \leq \gamma$  then  $\beta \alpha \leq \gamma \alpha$ .

Proof: Transfinite induction on  $\alpha$ .  $\square$ 

**Theorem 7.7.18** (Division Theorem). Let  $\alpha$  and  $\delta$  be ordinal numbers with  $\delta \neq 0$ . Then there exist unique ordinals  $\beta$  and  $\gamma$  with  $\gamma < \delta$  and

$$\alpha = \delta\beta + \gamma .$$

Proof:

- $\langle 1 \rangle 1$ . For any ordinal numbers  $\alpha$  and  $\delta$  with  $\delta \neq 0$ , there exist ordinals  $\beta$  and  $\gamma$  such that  $\gamma < \delta$  and  $\alpha = \delta \beta + \gamma$ 
  - $\langle 2 \rangle 1$ . Let:  $\alpha$  and  $\delta$  be ordinals with  $\delta \neq 0$
  - $\langle 2 \rangle 2$ . Let:  $\beta$  be the greatest ordinal such that  $\delta \beta \leq \alpha$

Proof: Proposition 7.6.6.

 $\langle 2 \rangle 3$ . There exists an ordinal  $\gamma$  such that  $\alpha = \delta \beta + \gamma$ 

PROOF: Subtraction Theorem

- $\langle 1 \rangle 2$ . For any ordinals  $\delta$ ,  $\beta$ ,  $\beta'$ ,  $\gamma$ ,  $\gamma'$ , if  $\delta \beta + \gamma = \delta \beta' + \gamma'$  and  $\delta \neq 0$  and  $\gamma, \gamma' < \delta$  then  $\beta = \beta'$  and  $\gamma = \gamma'$ 
  - $\langle 2 \rangle 1$ . Let:  $\delta$ ,  $\beta$ ,  $\beta'$ ,  $\gamma$ ,  $\gamma'$  be ordinals.
  - $\langle 2 \rangle 2$ . Assume:  $\delta \neq 0$  and  $\delta \beta + \gamma = \delta \beta' + \gamma'$
  - $\langle 2 \rangle 3. \ \beta = \beta'$ 
    - $\langle 3 \rangle 1. \ \beta \not< \beta'$

PROOF: If  $\beta < \beta'$  then

$$\delta\beta' + \gamma' \ge \delta\beta'$$

$$\ge \delta(\beta + 1)$$

$$= \delta\beta + \delta$$

$$> \delta\beta + \gamma$$

 $\langle 3 \rangle 2. \ \beta' \not < \beta$ 

PROOF: Similar.

 $\langle 2 \rangle 4. \ \gamma = \gamma'$ 

Proof: By Cancellation.

## 7.7.3 Exponentiation

**Definition 7.7.19.** Given ordinals  $\alpha$  and  $\beta$ , define the ordinal  $\alpha^{\beta}$  as follows:

$$\begin{array}{l} 0^{\alpha} := 0 & (\alpha > 0) \\ \alpha^{0} := 1 \\ \\ \alpha^{\beta^{+}} := \alpha^{\beta} \alpha & (\alpha > 0) \\ \\ \alpha^{\lambda} := \sup_{\beta < \lambda} \alpha^{\beta} & (\alpha > 0, \lambda \text{ a limit ordinal)} \end{array}$$

**Theorem 7.7.20.** Let  $\alpha$  be an ordinal  $\geq 2$ . The operation that maps  $\beta$  to  $\alpha^{\beta}$  is normal.

Proof:

 $\langle 1 \rangle 1$ . For  $\lambda$  a limit ordinal we have  $\alpha^{\lambda} = \sup_{\beta < \lambda} \alpha^{\beta}$  PROOF: By definition.

 $\langle 1 \rangle 2$ . For any ordinal  $\beta$  we have  $\alpha^{\beta} < \alpha^{\beta^+}$ 

PROOF: We have  $\alpha^{\beta^+} = \alpha^{\beta} \alpha > \alpha^{\beta}$  by Theorem 7.7.16 since  $\alpha > 1$  and  $\alpha^{\beta} \neq 0$ .

**Corollary 7.7.20.1.** For any ordinals  $\alpha$ ,  $\beta$ ,  $\gamma$ , if  $\alpha \geq 2$  then  $\beta < \gamma$  if and only if  $\alpha^{\beta} < \alpha^{\gamma}$ .

**Corollary 7.7.20.2** (Cancellation for Exponentiation). For any ordinals  $\alpha$ ,  $\beta$ ,  $\gamma$ , if  $\alpha \geq 2$  and  $\alpha^{\beta} = \alpha^{\gamma}$  then  $\beta = \gamma$ .

**Theorem 7.7.21.** For any ordinals  $\alpha$ ,  $\beta$  and  $\gamma$ , if  $\beta \leq \gamma$  then  $\beta^{\alpha} \leq \gamma^{\alpha}$ .

PROOF: Transfinite induction on  $\alpha$ .

**Theorem 7.7.22** (Logarithm Theorem). Let  $\alpha$  and  $\beta$  be ordinal numbers with  $\alpha \neq 0$  and  $\beta > 1$ . Then there exist unique ordinals  $\gamma$ ,  $\delta$  and  $\rho$  such that

$$\alpha = \beta^{\gamma} \delta + \rho, \qquad 0 \neq \delta < \beta, \qquad \rho < \beta^{\gamma}.$$

Proof:

 $\langle 1 \rangle 1$ . For any ordinals  $\alpha$  and  $\beta$  with  $\alpha \neq 0$  and  $\beta > 1$ , there exist ordinals  $\gamma$ ,  $\delta$ ,  $\rho$  such that

$$\alpha = \beta^{\gamma} \delta + \rho, \qquad 0 \neq \delta < \beta, \qquad \rho < \beta^{\gamma}.$$

- $\langle 2 \rangle 1$ . Let:  $\alpha$  and  $\beta$  be ordinals with  $\alpha \neq 0$  and  $\beta > 1$ .
- $\langle 2 \rangle 2$ . Let:  $\gamma$  be the greatest ordinal such that  $\beta^{\gamma} \leq \alpha$ .

Proof: Proposition 7.6.6.

- $\langle 2 \rangle$ 3. Let:  $\delta$  and  $\rho$  be the unique ordinals with  $\rho < \beta^{\gamma}$  such that  $\alpha = \beta^{\gamma} \delta + \rho$ . Proof: By the Division Theorem.
- $\langle 2 \rangle 4. \ \delta \neq 0$

PROOF: If  $\delta = 0$  then  $\alpha = \beta^{\gamma}0 + \rho = \rho < \beta^{\gamma} \le \alpha$  which is a contradiction.  $\langle 2 \rangle 5$ .  $\delta < \beta$ 

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PROOF: If  $\beta \leq \delta$  then  $\alpha \geq \beta^{\gamma} \delta \geq \beta^{\gamma} \beta = \beta^{\gamma+1}$ , contradicting the greatest-

- $\langle 1 \rangle 2$ . If  $\beta^{\gamma} \delta + \rho = \beta^{\gamma'} \delta' + \rho'$  with  $\beta > 1$ ,  $0 \neq \delta < \beta$ ,  $0 \neq \delta' < \beta$ ,  $\rho < \beta^{\gamma}$  and  $\rho' < \beta^{\gamma'}$ , then  $\gamma = \gamma'$ ,  $\delta = \delta'$  and  $\rho = \rho'$ .
  - $\langle 2 \rangle 1$ . Let:  $\alpha = \beta^{\gamma} \delta + \rho = \beta^{\gamma'} \delta' + \rho'$

  - $\begin{array}{l} \langle 2 \rangle 2. \ \beta^{\gamma} \leq \alpha < \beta^{\gamma+1} \\ \langle 2 \rangle 3. \ \beta^{\gamma'} \leq \alpha < \beta^{\gamma'+1} \\ \langle 2 \rangle 4. \ \beta^{\gamma} < \beta^{\gamma'+1} \ \text{and} \ \beta^{\gamma'} < \beta^{\gamma+1} \end{array}$
  - $\langle 2 \rangle 5$ .  $\gamma < \gamma' + 1$  and  $\gamma' < \gamma + 1$
  - $\langle 2 \rangle 6. \ \gamma = \gamma'$
  - $\langle 2 \rangle 7$ .  $\delta = \delta'$  and  $\rho = \rho'$

PROOF: By the Division Theorem.

**Theorem 7.7.23.** For any ordinal numbers  $\alpha$ ,  $\beta$ ,  $\gamma$ , we have

$$\alpha^{\beta+\gamma} = \alpha^{\beta}\alpha^{\gamma} .$$

Proof:

 $\langle 1 \rangle 1$ . Let:  $P[\gamma]$  be the property: for any ordinals  $\alpha$  and  $\beta$  we have  $\alpha^{\beta+\gamma}=\alpha^{\beta}\alpha^{\gamma}$  $\langle 1 \rangle 2$ . P[0]

Proof:

$$\alpha^{\beta+0} = \alpha^{\beta}$$
$$= \alpha^{\beta} 1$$
$$= \alpha^{\beta} \alpha^{0}$$

 $\langle 1 \rangle 3$ . For all  $\gamma$ , if  $P[\gamma]$  then  $P[\gamma + 1]$ 

Proof:

$$\begin{split} \alpha^{\beta+\gamma+1} &= \alpha^{\beta+\gamma}\alpha \\ &= \alpha^{\beta}\alpha^{\gamma}\alpha \\ &= \alpha^{\beta}\alpha^{\gamma+1} \end{split} \tag{induction hypothesis}$$

- $\langle 1 \rangle 4$ . For any limit ordinal  $\lambda$ , if  $\forall \gamma < \lambda . P[\gamma]$  then  $P[\lambda]$ .
  - $\langle 2 \rangle$ 1. Let:  $\lambda$  be a limit ordinal
  - $\langle 2 \rangle 2$ . Assume:  $\forall \gamma < \lambda . P[\gamma]$
  - $\langle 2 \rangle 3$ . Let:  $\alpha$  and  $\beta$  be any ordinals.
  - $\langle 2 \rangle 4$ . Case:  $\alpha = 0$

PROOF: We have  $\alpha^{\beta+\lambda} = \alpha^{\beta}\alpha^{\lambda} = 0$ .

 $\langle 2 \rangle$ 5. Case:  $\alpha = 1$ 

PROOF: We have  $\alpha^{\beta+\lambda} = \alpha^{\beta}\alpha^{\lambda} = 1$ .

 $\langle 2 \rangle 6$ . Case:  $\alpha > 1$ 

$$\begin{split} \alpha^{\beta+\lambda} &= \alpha^{\sup_{\gamma < \lambda} (\beta+\gamma)} \\ &= \sup_{\gamma < \lambda} \alpha^{\beta+\gamma} & \text{(Theorem 7.6.7)} \\ &= \sup_{\gamma < \lambda} \alpha^{\beta} \alpha^{\gamma} & \text{($\langle 2 \rangle 2$)} \\ &= \alpha^{\beta} \sup_{\gamma < \lambda} \alpha^{\gamma} & \text{(Theorem 7.6.7)} \\ &= \alpha^{\beta} \alpha^{\lambda} \end{split}$$

**Theorem 7.7.24.** For any ordinal numbers  $\alpha$ ,  $\beta$  and  $\gamma$ , we have

$$(\alpha^{\beta})^{\gamma} = \alpha^{\beta\gamma} .$$

Proof:

 $\langle 1 \rangle 1$ . Let:  $P[\gamma]$  be the property: For any ordinals  $\alpha$  and  $\beta$ , we have  $(\alpha^{\beta})^{\gamma} = \alpha^{\beta\gamma}$ 

 $\langle 1 \rangle 2$ . P[0] Proof:

$$(\alpha^{\beta})^0 = 1$$
$$= \alpha^{\beta 0}$$

 $\langle 1 \rangle 3. \ \forall \gamma \in \mathbf{On}.P[\gamma] \Rightarrow P[\gamma + 1]$ 

Proof:

$$(\alpha^{\beta})^{\gamma+1} = (\alpha^{\beta})^{\gamma} \alpha^{\beta}$$
$$= \alpha^{\beta\gamma} \alpha^{\beta}$$
$$= \alpha^{\beta\gamma+\beta}$$
$$= \alpha^{\beta(\gamma+1)}$$

 $\langle 1 \rangle 4$ . For any limit ordinal  $\lambda$ , if  $\forall \gamma < \lambda . P[\gamma]$  then  $P[\lambda]$ .

 $\langle 2 \rangle 1$ . Let:  $\lambda$  be a limit ordinal.

 $\langle 2 \rangle 2$ . Assume:  $\forall \gamma < \lambda . P[\gamma]$ 

 $\langle 2 \rangle 3$ . Let:  $\alpha$  and  $\beta$  be any ordinals.

 $\langle 2 \rangle 4$ . Case:  $\alpha = 0$  and  $\beta = 0$ 

PROOF:

$$(0^{\beta})^{\lambda} = 1^{\lambda}$$
$$= 1$$
$$= 0^{0}$$
$$= 0^{0\lambda}$$

 $\langle 2 \rangle$ 5. Case:  $\alpha = 0$  and  $\beta \neq 0$ Proof:  $(0^{\beta})^{\lambda} = 0^{\beta \lambda} = 0$ .

 $\langle 2 \rangle 6$ . Case:  $\alpha = 1$ 

PROOF:  $(1^{\beta})^{\lambda} = 1^{\beta\lambda} = 1$ 

 $\langle 2 \rangle 7$ . Case:  $\alpha > 1$ 

Proof:

$$(\alpha^{\beta})^{\lambda} = \sup_{\gamma < \lambda} (\alpha^{\beta})^{\gamma}$$
$$= \sup_{\gamma < \lambda} \alpha^{\beta\gamma}$$
$$= \alpha^{\sup_{\gamma < \lambda} \beta\gamma}$$
$$= \alpha^{\beta\lambda}$$

## 7.8 Sequences

i

**Definition 7.8.1** (Sequence). Given an ordinal  $\alpha$  and class  $\mathbf{A}$ , an  $\alpha$ -sequence in  $\mathbf{A}$  is a function  $a: \alpha \to \mathbf{A}$ . We write  $a_{\beta}$  for  $a(\beta)$ , and  $(a_{\beta})_{\beta < \alpha}$  for a.

**Definition 7.8.2** (Strictly Increasing). A sequence  $(a_{\beta})$  of ordinals is *strictly increasing* iff, whenever  $\beta < \gamma$ , then  $a_{\beta} < a_{\gamma}$ .

**Definition 7.8.3** (Subsequence). Let  $(a_{\beta})_{\beta<\gamma}$  be a sequence in **A**. A subsequence of  $(a_{\beta})$  is a sequence of the form  $(a_{\beta\xi})_{\xi<\delta}$  where  $(\beta\xi)_{\xi<\delta}$  is a strictly increasing sequence in  $\gamma$ .

**Definition 7.8.4** (Convergence). Let  $(a_{\beta})_{\beta < \gamma}$  be a sequence of ordinals and  $\lambda$  an ordinal. Then  $(a_{\beta})$  converges to the *limit*  $\lambda$  iff  $\lambda = \sup_{\beta < \gamma} a_{\beta}$ .

**Lemma 7.8.5.** Let  $(a_{\beta})_{\beta < \gamma}$  be a sequence of ordinals. Then there is a strictly increasing subsequence  $(a_{\beta \xi})_{\xi < \delta}$  such that  $\sup_{\xi < \delta} a_{\beta \xi} = \sup_{\beta < \gamma} a_{\beta}$ .

PROOF: Define  $\beta_{\xi}$  by transfinite recursion as follows.  $\beta_{\xi}$  is the least  $\beta$  such that  $a_{\beta} > a_{\beta_{\xi}}$  for all  $\zeta < \xi$  if there is such an  $a_{\beta}$ ; if not, the sequence ends.  $\square$ 

## 7.9 Strict Supremum

**Definition 7.9.1** (Strict Supremum). For any set S of ordinals, define the *strict* supremum of S, ssup S, to be the least ordinal greater than every member of S.

## Chapter 8

# Cardinal Numbers

## 8.1 Cardinal Numbers

**Definition 8.1.1** (Cardinality). For any set A, the *cardinality* or *cardinal number* |A| of A is the least ordinal equinumerous with A.

Let Card be the class of all cardinal numbers.

**Proposition 8.1.2.** For any sets A and B, we have  $A \approx B$  iff |A| = |B|.

Proof: Easy.  $\square$ 

**Definition 8.1.3** (Addition). Given cardinal numbers  $\kappa$  and  $\lambda$ , we define  $\kappa + \lambda$  to be  $|A \cup B|$  where A and B are disjoint sets of cardinality  $\kappa$  and  $\lambda$  respectively. We prove this is well-defined.

## Proof:

- $\langle 1 \rangle 1$ . Assume:  $A \approx A'$ ,  $B \approx B'$ , and  $A \cap B = A' \cap B' = \emptyset$
- $\langle 1 \rangle 2$ . Pick bijections  $f: A \approx A'$  and  $g: B \approx B'$
- $\langle 1 \rangle 3$ . The function  $A \cup B \to A' \cup B'$  that maps  $a \in A$  to f(a) and  $b \in B$  to g(b) is a bijection.

**Proposition 8.1.4.** For any cardinal number  $\kappa$ , we have  $\kappa + 0 = \kappa$ .

PROOF: Let A and B be disjoint sets of cardinality  $\kappa$  and A. Then  $A = \emptyset$  so  $A \cup B = A$  and so  $|A \cup B| = \kappa$ .  $\square$ 

**Theorem 8.1.5** (Associative Law for Addition). For any cardinal numbers  $\kappa$ ,  $\lambda$ ,  $\mu$  we have  $\kappa + (\lambda + \mu) = (\kappa + \lambda) + \mu$ .

PROOF: Since  $A \cup (B \cup C) = (A \cup B) \cup C$ .  $\square$ 

**Proposition 8.1.6.** For any cardinal numbers  $\kappa$  and  $\lambda$  we have  $\kappa + \lambda = \lambda + \kappa$ .

PROOF: Since  $A \cup B = B \cup A$ .  $\square$ 

**Definition 8.1.7** (Multiplication). For  $\kappa$  and  $\lambda$  cardinal numbers, we define  $\kappa\lambda$  to be the cardinal number of  $A\times B$ , where  $|A|=\kappa$  and  $|B|=\lambda$ .

We prove this is well-defined.

PROOF: If  $f: A \approx A'$  and  $g: B \approx B'$  then the function that maps (a,b) to (f(a),g(b)) is a bijection  $A \times B \approx A' \times B'$ .  $\square$ 

**Proposition 8.1.8.** For any cardinal number  $\kappa$  we have  $\kappa \cdot 0 = 0$ .

PROOF: Since  $A \times \emptyset = \emptyset$ .  $\square$ 

**Proposition 8.1.9.** For any cardinal number  $\kappa$  we have  $\kappa \cdot 1 = \kappa$ .

PROOF: The function that maps (a, e) to a is a bijection  $A \times \{e\} \approx A$ .  $\square$ 

**Theorem 8.1.10** (Distributive Law). For any cardinal numbers  $\kappa$ ,  $\lambda$  and  $\mu$ , we have  $\kappa(\lambda + \mu) = \kappa\lambda + \kappa\mu$ .

PROOF: Since  $A \times (B \cup C) = (A \times B) \cup (A \times C)$ .  $\square$ 

**Theorem 8.1.11** (Associative Law for Multiplication). For any cardinal numbers  $\kappa$ ,  $\lambda$  and  $\mu$ , we have  $\kappa(\lambda\mu) = (\kappa\lambda)\mu$ .

PROOF: Since  $A \times (B \times C) \approx (A \times B) \times C$ .  $\square$ 

**Theorem 8.1.12** (Commutative Law for Multiplication). For any cardinal numbers  $\kappa$  and  $\lambda$ , we have  $\kappa\lambda = \lambda\kappa$ .

PROOF: Since  $A \times B \approx B \times A$ .  $\square$ 

**Theorem 8.1.13.** For any cardinal numbers  $\kappa$  and  $\lambda$ , if  $\kappa\lambda = 0$  then  $\kappa = 0$  or  $\lambda = 0$ .

PROOF: if  $A \times B = \emptyset$  then  $A = \emptyset$  or  $B = \emptyset$ .  $\square$ 

**Definition 8.1.14** (Exponentiation). Given cardinal numbers  $\kappa$  and  $\lambda$ , we define  $\kappa^{\lambda}$  to be  $|A^{B}|$ , where  $|A| = \kappa$  and  $|B| = \lambda$ .

We prove this is well-defined.

PROOF:If  $f: A \approx A'$  and  $g: B \approx B'$ , then the function that maps  $h: B \to A$  to  $f \circ h \circ g^{-1}$  is a bijection  $A^B \approx A'^{B'}$ .  $\square$ 

**Proposition 8.1.15.** For any cardinal numbers  $\kappa$ ,  $\lambda$  and  $\mu$ ,

$$\kappa^{\lambda+\mu} = (\kappa^{\lambda})^{\mu}$$

PROOF: The function that maps  $f: A \times B \to C$  to  $\lambda a \in A.\lambda b \in B.f(a,b)$  is a bijection  $A^{B \times C} \approx (A^B)^C$ .  $\square$ 

**Proposition 8.1.16.** For any cardinal numbers  $\kappa$ ,  $\lambda$  and  $\mu$ ,

$$(\kappa\lambda)^{\mu} = \kappa^{\mu}\lambda^{\mu}$$
.

PROOF: The function  $f: A^C \times B^C \to (A \times B)^C$  with f(g,h)(c) = (g(c),h(c)) is a bijection.  $\square$ 

**Proposition 8.1.17.** For any cardinal numbers  $\kappa$ ,  $\lambda$  and  $\mu$ , we have

$$\kappa^{\lambda+\mu} = \kappa^{\lambda} \kappa^{\mu}$$
.

PROOF: If  $B \cap C = \emptyset$ , then  $f: A^B \times A^C \to A^{B \cup C}$  given by f(g,h)(b) = g(b)and f(g,h)(c) = h(c) is a bijection.  $\square$ 

**Proposition 8.1.18.** For any cardinal number  $\kappa$ , we have  $\kappa^0 = 1$ .

PROOF: For any set A, we have  $A^{\emptyset} = \{\emptyset\}$ .  $\square$ 

**Proposition 8.1.19.** For any cardinal number  $\kappa$ , we have  $\kappa^1 = \kappa$ .

PROOF: For any sets A and B, if  $B = \{b\}$  then the function  $f: A \to A^B$  with f(a)(b) = a is a bijection.  $\square$ 

**Proposition 8.1.20.** For any non-zero cardinal number  $\kappa$  we have  $0^{\kappa} = 0$ .

PROOF: If A is nonempty then there is no function  $A \to \emptyset$ .  $\square$ 

**Proposition 8.1.21.** For any set A we have  $|\mathcal{P}A| = 2^{|A|}$ .

PROOF: The function  $f: \mathcal{P}A \to 2^A$  where f(X)(a) = 0 if  $a \notin X$  and f(X)(a) = 01 if  $a \in X$ .  $\square$ 

**Theorem 8.1.22** (König). Let I be a set. Let  $\{A_i\}_{i\in I}$  and  $\{B_i\}_{i\in I}$  be families of sets. Assume that  $\forall i \in I. |A_i| < |B_i|$ . Then  $|\bigcup_{i \in I} A_i| < |\prod_{i \in I} B_i|$ .

#### Proof:

$$\begin{array}{l} \langle 1 \rangle 1. \text{ For all } i \in I, \text{ choose an injection } f_i : A_i \rightarrowtail B_i \\ \langle 1 \rangle 2. \text{ For all } i \in I, \text{ choose } b_i \in B_i - f_i(A_i) \\ \langle 1 \rangle 3. \left| \bigcup_{i \in I} A_i \right| \leq \left| \prod_{i \in I} B_i \right| \\ \langle 2 \rangle 1. \text{ Define } g : \bigcup_{i \in I} A_i \to \prod_{i \in I} B_i \text{ by} \\ g(i,a)(j) = \begin{cases} f_i(a) & \text{if } i = j \\ b_j & \text{otherwise} \end{cases} \\ \langle 2 \rangle 2. g \text{ is injective.} \\ \langle 1 \rangle 4. \left| \bigcup_{i \in I} A_i \right| \neq \left| \prod_{i \in I} B_i \right| \\ \end{array}$$

- $\langle 1 \rangle 4$ .  $\left| \bigcup_{i \in I} A_i \right| \neq \left| \prod_{i \in I} B_i \right|$  $\langle 2 \rangle 1$ . Let:  $h : \bigcup_{i \in I} A_i \to \prod_{i \in I} B_i$ Prove: h is not surjective.
  - $\langle 2 \rangle 2$ . For  $i \in I$ , Pick  $c_i \in B_i \{h(i, a)(i) \mid i \in I\}$
- $\langle 2 \rangle 3. \ c \in \prod_{i \in I} B_i$
- $\langle 2 \rangle 4$ .  $c \notin \operatorname{ran} h$

Corollary 8.1.22.1. For any cardinal number  $\kappa$  we have  $\kappa < 2^{\kappa}$ .

## 8.2 Ordering on Cardinal Numbers

**Definition 8.2.1.** Given cardinal numbers  $\kappa$  and  $\lambda$ , we have  $\kappa \leq \lambda$  iff  $A \leq B$ , where  $|A| = \kappa$  and  $|B| = \lambda$ .

```
Proof:
\langle 1 \rangle 1. Let: |A| = \kappa and |B| = \lambda
(1)2. Pick bijections f: A \approx \kappa and g: B \approx \lambda
\langle 1 \rangle 3. If \kappa \leq \lambda then A \preccurlyeq B
    PROOF: Let i: \kappa \hookrightarrow \lambda be the inclusion. Then g^{-1} \circ i \circ f is an injection A \to B.
\langle 1 \rangle 4. If A \leq B then \kappa \leq \lambda
    \langle 2 \rangle 1. Assume: A \leq B
    \langle 2 \rangle 2. Pick an injection h: A \rightarrow B
    \langle 2 \rangle 3. g(h(A)) \subseteq B is well-ordered by \in
    \langle 2 \rangle 4. Let: \gamma be the ordinal number of (g(h(A)), \in)
    \langle 2 \rangle 5. \ \gamma \leq \lambda
       Proof: Proposition 7.3.13.
    \langle 2 \rangle 6. \ \kappa \leq \gamma
       PROOF: By the leastness of \kappa, since A is equinumerous with \gamma.
    \langle 2 \rangle 7. \ \kappa \leq \lambda
П
```

Corollary 8.2.1.1. There is no largest cardinal number.

**Proposition 8.2.2.** For any cardinal numbers  $\kappa$ ,  $\lambda$ ,  $\mu$ , if  $\kappa \leq \lambda$  then  $\kappa + \mu \leq \lambda + \mu$ .

PROOF: If  $f: A \to B$  is injective, and  $A \cap C = B \cap C = \emptyset$ , then the function  $A \cup C \to B \cup C$  that maps a to f(a) and maps c to c is an injection.  $\square$ 

**Proposition 8.2.3.** For any cardinal numbers  $\kappa$ ,  $\lambda$ ,  $\mu$ , if  $\kappa \leq \lambda$  then  $\kappa \mu \leq \lambda \mu$ .

PROOF: If  $f: A \to B$  is injective, then the function  $A \times C \to B \times C$  that maps (a,c) to (f(a),c) is injective.  $\square$ 

**Proposition 8.2.4.** For any cardinal numbers  $\kappa$ ,  $\lambda$ ,  $\mu$ , if  $\kappa \leq \lambda$  then  $\kappa^{\mu} \leq \lambda^{\mu}$ .

PROOF: Given an injection  $f:A\to B$ , the function that maps  $A^C\to B^C$  that maps g to  $f\circ g$  is an injection.  $\square$ 

**Proposition 8.2.5.** For any cardinal numbers  $\kappa$ ,  $\lambda$ ,  $\mu$ , if  $\kappa \leq \lambda$  and  $\mu$  and  $\kappa$  are not both 0, then  $\mu^{\kappa} \leq \mu^{\lambda}$ .

Proof:

- $\langle 1 \rangle 1$ . Let: A, B and C be sets with A and C not both empty.
- $\langle 1 \rangle 2$ . Let:  $f: A \to B$  be an injection. Prove:  $C^A \preccurlyeq C^B$

 $\langle 1 \rangle 3$ . Case:  $C = \emptyset$ 

PROOF: Then  $A \neq \emptyset$  so  $C^A = \emptyset \preccurlyeq C^B$ .

 $\langle 1 \rangle 4$ . Case:  $C \neq \emptyset$ 

 $\langle 2 \rangle 1$ . Pick  $c \in C$ 

 $\langle 2 \rangle 2$ . Let:  $g: C^A \to C^B$  be the function g(h)(f(a)) = h(a), g(h)(b) = c if

 $\langle 2 \rangle 3$ . g is an injection.

**Proposition 8.2.6.** Let A be a set such that  $\forall X \in A | X | \leq \kappa$ . Then

$$\left|\bigcup \mathcal{A}\right| \leq |\mathcal{A}|\kappa \ .$$

## Proof:

 $\langle 1 \rangle 1$ . For  $X \in \mathcal{A}$ , choose a surjection  $f_X : \kappa \to X$ .

 $\langle 1 \rangle 2$ . Define  $g: \mathcal{A} \times \kappa \to \bigcup \mathcal{A}$  by  $g(X, \alpha) = f_X(\alpha)$ 

 $\langle 1 \rangle 3$ . g is surjective.

**Lemma 8.2.7.** The union of a set of cardinal numbers is a cardinal number.

 $\langle 1 \rangle 1$ . Let: A be a set of cardinal numbers.

PROVE:  $\bigcup A$  is the smallest ordinal equinumerous with  $\bigcup A$ 

 $\langle 1 \rangle 2$ . Let:  $\alpha < \bigcup A$ 

Prove:  $\alpha \not\approx \bigcup A$ 

 $\langle 1 \rangle 3$ . Pick  $\kappa \in A$  such that  $\alpha < \kappa$ 

 $\langle 1 \rangle 4$ .  $\alpha \prec \kappa$ 

 $\langle 1 \rangle 5. \ \stackrel{\backsim}{\alpha} \stackrel{\backsim}{\prec} \stackrel{\kappa}{\bigcup} A$ 

# Chapter 9

# **Natural Numbers**

## 9.1 Inductive Sets

**Definition 9.1.1** (Inductive). A set I is *inductive* iff  $0 \in I$  and  $\forall x \in I.x^+ \in I$ .

**Definition 9.1.2** (Natural Number). A *natural number* is a set that belongs to every inductive set.

**Theorem 9.1.3.** The class  $\mathbb{N}$  of natural numbers is a set.

```
PROOF: \langle 1 \rangle 1. PICK an inductive set I. PROOF: Axiom of Infinity. \langle 1 \rangle 2. \mathbb{N} \subseteq I
```

**Theorem 9.1.4.**  $\mathbb{N}$  is inductive, and is a subset of every other inductive set.

```
PROOF: \langle 1 \rangle 1. \mathbb{N} is inductive. \langle 2 \rangle 1. 0 \in \mathbb{N}
PROOF: Since 0 is a member of every inductive set. \langle 2 \rangle 2. \forall n \in \mathbb{N}. n^+ \in \mathbb{N}
\langle 3 \rangle 1. Let: n \in \mathbb{N}
\langle 3 \rangle 2. Let: I be any inductive set.
PROVE: n^+ \in I
\langle 3 \rangle 3. n \in I
PROOF: \langle 3 \rangle 1, \langle 3 \rangle 2
\langle 3 \rangle 4. n^+ \in I
PROOF: \langle 3 \rangle 2, \langle 3 \rangle 3
\langle 1 \rangle 2. \mathbb{N} is a subset of every inductive set.
PROOF: Immediate from definitions.
```

Corollary 9.1.4.1 (Induction Principle for  $\mathbb{N}$ ). Any inductive subset of  $\mathbb{N}$  coincides with  $\mathbb{N}$ .

**Theorem 9.1.5.** Every natural number except 0 is the successor of some natural number.

Proof: Trivially by induction.

Proposition 9.1.6. Every natural number is an ordinal.

Proof: By induction.  $\square$ 

**Proposition 9.1.7.**  $\mathbb{N}$  is a transitive set.

#### Proof:

- $\langle 1 \rangle 1. \ 0 \subseteq \mathbb{N}$
- $\langle 1 \rangle 2. \ \forall n \in \mathbb{N}. n \subseteq \mathbb{N} \Rightarrow n^+ \subseteq \mathbb{N}$
- $\langle 1 \rangle 3. \ \forall n \in \mathbb{N}. n \subseteq \mathbb{N}$

PROOF: From  $\langle 1 \rangle 1$  and  $\langle 1 \rangle 2$  by induction.

Corollary 9.1.7.1.  $\mathbb{N}$  is an ordinal.

**Definition 9.1.8.** We define  $\omega = \mathbb{N}$ .

**Proposition 9.1.9** (Dependent Choice). Let A be a nonempty set and R a relation on A such that  $\forall x \in A. \exists y \in A. (y,x) \in R$ . Then there exists a function  $f: \mathbb{N} \to A$  such that  $\forall n \in \mathbb{N}. (f(n+1), f(n)) \in R$ .

### Proof:

- $\langle 1 \rangle 1$ . PICK a choice function F for A.
- $\langle 1 \rangle 2$ . Pick  $a \in A$
- $\begin{array}{l} \langle 1 \rangle 3. \text{ Define } f: \mathbb{N} \to A \text{ by } f(0) = a \text{ and } f(n+1) = F(\{y \in A \mid (y,f(n)) \in R\}). \end{array}$

**Theorem Schema 9.1.10.** For any classes A and R, the following is a theorem:

Assume **R** is a relation on **A** and, for all  $a \in \mathbf{A}$ , the class  $\{x \in \mathbf{A} \mid x\mathbf{R}a\}$  is a set. Then **R** is well founded if and only if there does not exist a function  $f: \mathbb{N} \to \mathbf{A}$  such that  $\forall n \in \mathbb{N}. f(n+1)\mathbf{R}f(n)$ .

#### Proof:

 $\langle 1 \rangle 1$ . If there exists a function  $f : \mathbb{N} \to \mathbf{A}$  such that  $\forall n \in \mathbb{N}. f(n+1)\mathbf{R}f(n)$  then  $\mathbf{R}$  is not well founded.

PROOF:  $f(\mathbb{N})$  is a nonempty subset of **A** with no **R**-minimal element.

- $\langle 1 \rangle$ 2. If **R** is not well founded then there exists a function  $f : \mathbb{N} \to \mathbf{A}$  such that  $\forall n \in \mathbb{N}. f(n+1)\mathbf{R}f(n)$ .
  - $\langle 2 \rangle 1$ . Assume: **R** is not well founded.
  - $\langle 2 \rangle 2$ . PICK a nonempty subset  $B \subseteq \mathbf{A}$  that has no **R**-minimal element.
  - $\langle 2 \rangle 3. \ \forall x \in B. \exists y \in B. y \mathbf{R} x$

```
\langle 2 \rangle 4. Choose a function g: B \to B such that \forall x \in B.g(x)\mathbf{R}x \langle 2 \rangle 5. PICK b \in B \langle 2 \rangle 6. Define f: \mathbb{N} \to \mathbf{A} recursively by f(0) = b and \forall n \in \mathbb{N}.f(n+1) = g(f(n)) \langle 2 \rangle 7. \forall n \in \mathbb{N}.f(n+1)\mathbf{R}f(n)
```

## 9.2 Cardinality

PROOF:  $\langle 2 \rangle 1$  $\langle 3 \rangle 5$ . ran f = n + 1

 $\langle 1 \rangle 4$ . For every natural number n, P(n).

**Definition 9.2.1** (Finite). A set is *finite* iff it is equinumerous to some natural number; otherwise it is *infinite*.

**Theorem 9.2.2** (Pigeonhole Principle). No natural number is equinumerous to a proper subset of itself.

```
Proof:
\langle 1 \rangle 1. Let: P(n) be the property: any one-to-one function n \to n is surjective.
\langle 1 \rangle 2. P(0)
   PROOF: The only function 0 \to 0 is injective.
\langle 1 \rangle 3. For every natural number n, if P(n) then P(n+1).
    \langle 2 \rangle 1. Assume: P(n)
    \langle 2 \rangle 2. Let: f be a one-to-one function n+1 \to n+1
   \langle 2 \rangle 3. f \upharpoonright n is a one-to-one function n \to n+1
   \langle 2 \rangle 4. Case: n \notin ranf
       \langle 3 \rangle 1. \ f \upharpoonright n : n \to n
       \langle 3 \rangle 2. ran(f \upharpoonright n) = n
       \langle 3 \rangle 3. \ f(n) = n
          Proof: \langle 2 \rangle 1.
       \langle 3 \rangle 4. ran f = n + 1
   \langle 2 \rangle 5. Case: n \in \operatorname{ran} f
       \langle 3 \rangle 1. PICK p \in n such that f(p) = n
       \langle 3 \rangle 2. Let: \hat{f}: n \to n be the function
                                                \hat{f}(p) = f(n)
                                                \hat{f}(x) = f(x)
                                                                                  (x \neq p)
       \langle 3 \rangle 3. \hat{f} is one-to-one
       \langle 3 \rangle 4. ran \hat{f} = n
```

Corollary 9.2.2.1. No finite set is equinumerous to a proper subset of itself.

Corollary 9.2.2.2. Every natural number is a cardinal number.

PROOF: For any natural number n, we have that n is the least ordinal such that  $n \approx n$ .  $\square$ 

Corollary 9.2.2.3.  $\mathbb{N}$  is a cardinal number.

Corollary 9.2.2.4.  $\mathbb{N}$  is infinite.

PROOF: The function that maps n to n+1 is a bijection between  $\mathbb{N}$  and  $\mathbb{N}-\{0\}$ .

**Corollary 9.2.2.5.** If C is a proper subset of a natural number n, then there exists m < n such that  $C \approx m$ .

Proof: By Proposition 7.3.13.  $\square$ 

Corollary 9.2.2.6. Any subset of a finite set is finite.

**Proposition 9.2.3.** For any natural numbers m and n we have m+n (cardinal addition) is a natural number.

PROOF: Induction on n.  $\square$ 

Corollary 9.2.3.1. The union of two finite sets is finite.

Corollary 9.2.3.2. The union of a finite set of finite sets is finite.

PROOF: By induction on the number of elements.  $\square$ 

**Proposition 9.2.4.** For natural numbers m and n, the cardinal sum m + n is equal to the ordinal sum m + n.

Proof: Induction on n.

**Proposition 9.2.5.** For any natural numbers m and n, we have mn (cardinal multiplication) is a natural number.

**Corollary 9.2.5.1.** If A and B are finite sets then  $A \times B$  is finite.

**Proposition 9.2.6.** For natural numbers m and n, the cardinal product mn is equal to the ordinal product mn.

PROOF: Induction on n.

**Proposition 9.2.7.** For any natural numbers m and n we have  $m^n$  (cardinal exponentiation) is a natural number.

PROOF: Induction on n.

Corollary 9.2.7.1. If A and B are finite sets then  $A^B$  are finite.

**Proposition 9.2.8.** For natural numbers m and n, the cardinal exponentiation  $m^n$  and the ordinal exponentiation  $m^n$  agree.

PROOF: Induction on n.  $\square$ 

Proposition 9.2.9.  $\mathbb{N}^2 \approx \mathbb{N}$ 

PROOF: The function  $J: \mathbb{N}^2 \to \mathbb{N}$  defined by  $J(m,n) = ((m+n)^2 + 3m + n)/2$  is a bijection.  $\square$ 

**Proposition 9.2.10.** For any infinite cardinal  $\kappa$  we have  $\aleph_0 \leq \kappa$ .

#### Proof:

 $\langle 1 \rangle 1$ . Let: A be an infinite set.

Prove:  $\mathbb{N} \preceq A$ 

 $\langle 1 \rangle 2$ . PICK a choice function F for A.

 $\langle 1 \rangle 3$ . Define  $h : \mathbb{N} \to \{ X \in \mathcal{P}A \mid X \text{ is finite} \}$  by  $h(0) = \emptyset$ 

$$h(n+1) = h(n) \cup \{F(A - \{h(m) \mid m < n\})\}$$

 $\langle 1 \rangle 4$ . Define  $g : \mathbb{N} \to A$  by  $g(n) = F(A - \{h(m) \mid m < n\})$ 

 $\langle 1 \rangle 5$ . g is injective.

PROOF: If m < n then  $g(m) \neq g(n)$ .

**Theorem Schema 9.2.11** (König's Lemma). For any classes A and R, the following is a theorem:

Assume **R** is a well founded relation on **A** such that, for all  $y \in \mathbf{A}$ , the class  $\{x \in \mathbf{A} \mid x\mathbf{R}y\}$  is a finite set. Let  $\mathbf{R}^t$  be the transitive closure of **R**. Then, for all  $y \in \mathbf{A}$ , the class  $\{x \in \mathbf{A} \mid x\mathbf{R}^ty\}$  is a finite set.

## Proof:

- $\langle 1 \rangle 1$ . Let:  $y \in \mathbf{A}$
- $\langle 1 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall x \mathbf{R} y . \{z \in \mathbf{A} \mid z \mathbf{R}^t x\}$  is a finite set.

 $\langle 1 \rangle 3. \{ x \mid x \mathbf{R}^t y \} = \bigcup_{x \mathbf{R} y} (\{x\} \cup \{z \mid z \mathbf{R}^t x \})$ 

 $\langle 1 \rangle 4$ .  $\{x \mid x \mathbf{R}^t y\}$  is finite.

Proof: Corollary 9.2.3.2.

## 9.3 Countable Sets

**Definition 9.3.1** (Countable). A set A is countable iff  $|A| \leq \aleph_0$ .

**Theorem 9.3.2.** The union of a countable set of countable sets is countable.

Proof: Proposition 8.2.6.  $\square$ 

## 9.4 Arithmetic

**Definition 9.4.1** (Even). A natural number n is *even* iff there exists  $m \in \mathbb{N}$  such that n = 2m.

**Definition 9.4.2** (Odd). A natural number n is odd iff there exists  $p \in \mathbb{N}$  such that n = 2p + 1.

```
Proposition 9.4.3. Every natural number is either even or odd.
Proof:
\langle 1 \rangle 1. 0 is even.
  Proof: 0 = 2 \times 0.
\langle 1 \rangle 2. For any natural number n, if n is either even or odd then n^+ is either even
  Proof:
   \langle 2 \rangle 1. Let: n \in \mathbb{N}
   \langle 2 \rangle 2. If n is even then n^+ is odd.
     PROOF: If n = 2p then n^+ = 2p + 1.
   \langle 2 \rangle 3. If n is odd then n^+ is even.
     PROOF: If n = 2p + 1 then n^{+} = 2(p + 1).
Proposition 9.4.4. No natural number is both even and odd.
Proof:
\langle 1 \rangle 1. 0 is not odd.
  PROOF: For any p we have 2p + 1 = (2p)^+ \neq 0.
\langle 1 \rangle 2. For any natural number n, if n is not both even and odd, then n^+ is not
```

- both even and odd.  $\langle 2 \rangle 1$ . Let: n be a natural number.
- $\langle 2 \rangle 2$ . If  $n^+$  is even then n is odd.
  - $\langle 3 \rangle 1$ . Assume:  $n^+$  is even.
  - $\langle 3 \rangle 2$ . PICK p such that  $n^+ = 2p$
  - $\langle 3 \rangle 3. \ p \neq 0$

PROOF: Since  $n^+ \neq 0$ .

 $\langle 3 \rangle 4$ . Pick q such that  $p = q^+$ PROOF: Theorem 9.1.5.

 $\langle 3 \rangle 5. \ n^+ = 2q + 2$ 

Proof:  $\langle 3 \rangle 2$ ,  $\langle 3 \rangle 4$ .

 $\langle 3 \rangle 6. \ \ n = 2q + 1$ 

Proof: Proposition 7.4.7,  $\langle 3 \rangle 5$ 

- $\langle 3 \rangle 7$ . *n* is odd.
- $\langle 2 \rangle 3$ . If  $n^+$  is odd then n is even.
  - $\langle 3 \rangle 1$ . Assume:  $n^+$  is odd.
  - $\langle 3 \rangle 2$ . Pick p such that  $n^+ = 2p + 1$
  - $\langle 3 \rangle 3$ . n = 2p

Proof: Proposition 7.4.7,  $\langle 3 \rangle 2$ 

 $\langle 3 \rangle 4$ . *n* is even.

**Proposition 9.4.5.** Let m, n, p, q be natural numbers. Assume m+n=p+q. Then m < p if and only if q < n.

#### Proof:

 $\langle 1 \rangle 1$ . If m < p then q < n.

PROOF: If m < p and  $n \le q$  then  $m + n . <math>\langle 1 \rangle 2$ . If q < n then m < p. PROOF: Similar.  $\Box$ 

**Proposition 9.4.6.** Let m, n, p and q be natural numbers. Assume n < m and q < p. Then

$$mq + np < mp + nq$$
.

Proof:

 $\langle 1 \rangle 1$ . PICK positive natural numbers a and b such that m=n+a and p=q+b.  $\langle 1 \rangle 2$ . mp+nq>mq+np

PROOF:

$$mp + nq = (n+a)(q+b) + nq$$

$$= 2nq + nb + aq + ab$$

$$mq + np = (n+a)q + n(q+b)$$

$$= 2nq + aq + nb$$

$$\therefore mp + nq = mq + np + ab$$

$$> mq + np$$

## 9.5 Sequences

**Definition 9.5.1** (Sequence). Let A be a set. A *finite sequence* in A is a function  $a:n\to A$  for some natural number n; we write it as  $(a(0),a(1),\ldots,a(n-1))$ . An *(infinite) sequence* in A is a function  $\mathbb{N}\to A$ .

We write  $A^*$  for the set of all finite sequences in A.

**Proposition 9.5.2.** If A is countable then  $A^*$  is countable.

PROOF: For any n, the set  $A^n$  is countable, and  $A^*$  is equinumerous with  $\bigcup_n A^n$ .  $\sqcap$ 

## 9.6 Transitive Closure of a Set

**Proposition 9.6.1.** For any set A, there exists a unique transitive set C such that:

- $A \subseteq C$
- For any transitive set X, if  $A \subseteq X$  then  $C \subseteq X$

Proof:

 $\langle 1 \rangle 1.$  Define a function  $F: \mathbb{N} \to \mathbf{V}$  by F(0) = A  $F(n+1) = A \cup \bigcup (F(0) \cup \cdots \cup F(n))$ 

```
\langle 1 \rangle 2. For all n \in \mathbb{N} and a \in F(n) we have a \subseteq F(n+1)
   PROOF: a \in F(0) \cup \cdots \cup F(n) so a \subseteq \bigcup (F(0) \cup \cdots \cup F(n)) \subseteq F(n+1).
\langle 1 \rangle 3. Let: C = \bigcup_{n \in \mathbb{N}} F(n)
\langle 1 \rangle 4. C is transitive.
    \langle 2 \rangle 1. Let: x \in y \in C
    \langle 2 \rangle 2. Pick n \in \mathbb{N} such that y \in F(n)
    \langle 2 \rangle 3. \ y \subseteq F(n+1)
       Proof: \langle 1 \rangle 2
    \langle 2 \rangle 4. \ x \in F(n+1)
    \langle 2 \rangle 5. \ x \in C
\langle 1 \rangle 5. A \subseteq C
   PROOF: Since F(0) = A.
\langle 1 \rangle 6. For any transitive set X, if A \subseteq X then C \subseteq X
    \langle 2 \rangle 1. Let: X be a transitive set
    \langle 2 \rangle 2. Assume: A \subseteq X
    \langle 2 \rangle 3. For all n \in \mathbb{N} we have F(n) \subseteq X.
       \langle 3 \rangle 1. \ F(0) \subseteq X
           Proof: \langle 2 \rangle 2
       \langle 3 \rangle 2. For all n \in \mathbb{N}, if F(n) \subseteq X, then F(n+1) \subseteq X.
           \langle 4 \rangle 1. Let: n \in \mathbb{N}
           \langle 4 \rangle 2. Assume: \forall m < n.F(m) \subseteq X
           \langle 4 \rangle 3. \ F(0) \cup \cdots \cup F(n) \subseteq X
           \langle 4 \rangle 4. \bigcup (F(0) \cup \cdots \cup F(n)) \subseteq X
              Proof: Since X is transitive.
           \langle 4 \rangle 5. F(n+1) \subseteq X
    \langle 2 \rangle 4. \ C \subseteq X
\langle 1 \rangle 7. Let D be a transitive set such that A \subseteq D and, for any transitive set X,
         if A \subseteq X then D \subseteq X. Then D = C.
   PROOF: We have C \subseteq D and D \subseteq C.
Proposition 9.6.2 (ZF). For any class A, the relation \mathbf{E} = \{(x,y) \in \mathbf{A}^2 \mid x \in
y} is well founded.
Proof:
\langle 1 \rangle 1. Every nonempty set has an E-minimal element.
   Proof: Axiom of Regularity.
\langle 1 \rangle 2. For every set x, there exists a set u such that x \subseteq u and, for all w, y, if
         y \in u and w\mathbf{E}y then w \in u.
   PROOF: Take u to be the transitive closure of x.
```

## 9.7 The Veblen Fixed Point Theorem

**Theorem Schema 9.7.1** (Veblen Fixed Point Theorem). For any class  $\mathbf{T}$ , the following is a theorem:

Assume T is a normal ordinal operation. For every ordinal  $\beta$ , there exists  $\gamma \geq \beta$  such that  $\mathbf{T}(\gamma) = \gamma$ .

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $\beta$  be an ordinal.
- $\langle 1 \rangle 2$ . Assume: w.l.o.g.  $\beta < \mathbf{T}(\beta)$

PROOF: We have  $\beta \leq \mathbf{T}(\beta)$  by Proposition 7.6.5, and if  $\beta = \mathbf{T}(\beta)$  we take  $\gamma := \beta$ .

 $\langle 1 \rangle 3$ . Define  $f : \mathbb{N} \to \mathbf{On}$  by recursion thus:

$$f(0) = \beta$$
$$f(n^+) = \mathbf{T}(f(n))$$

$$\langle 1 \rangle 4$$
. Let:  $\gamma = \sup_{n \in \mathbb{N}} f(n)$ 

 $\langle 1 \rangle 5. \ \beta \leq \gamma$ 

PROOF: Since  $\beta = f(0)$ .

$$\langle 1 \rangle 6. \ \mathbf{T}(\gamma) = \gamma$$

$$\langle 2 \rangle 1. \ \mathbf{T}(\gamma) \leq \gamma$$

Proof:

$$\mathbf{T}(\gamma) = \sup_{n \in \mathbb{N}} \mathbf{T}(f(n))$$
 (Theorem 7.6.7)  
$$= \sup_{n \in \mathbb{N}} f(n^{+})$$
 (\langle 1\rangle 3)  
$$\leq \sup_{n \in \mathbb{N}} f(n)$$
  
$$= \gamma$$

 $\langle 2 \rangle 2. \ \gamma \leq \mathbf{T}(\gamma)$ 

Proof:Proposition 7.6.5.

**Definition 9.7.2** (Derived Operation). Let **T** be a normal ordinal operation. The derived operation  $T': On \to V$  is the unique order isomorphism between On and the fixed points of T.

**Proposition Schema 9.7.3.** For any class **T**, the following is a theorem: If **T** is a normal ordinal operation, then the derived operation is normal.

## Proof:

 $\langle 1 \rangle 1$ . For any set S of fixed points of **T**, we have  $\bigcup S$  is a fixed point of **T**  $\langle 2 \rangle 1$ . Let: S be a set of fixed points of **T**.

$$\langle 2 \rangle 2$$
.  $\mathbf{T}(\sup S) = \sup S$ 

Proof:

$$\mathbf{T}(\sup S) = \sup_{\alpha \in S} \mathbf{T}(\alpha)$$
 (Theorem 7.6.7)  
$$= \sup_{\alpha \in S} \alpha$$
 ( $\langle 2 \rangle 1$ )  
$$= \sup S$$

 $\langle 1 \rangle 2$ . Q.E.D.

Proof: Proposition 7.6.8.

## 9.8 Cantor Normal Form

**Theorem 9.8.1.** For any ordinal  $\alpha$ , there exist a unique sequence of nonzero natural numbers  $(n_1, \ldots, n_k)$  and sequence of ordinals  $(\gamma_1, \ldots, \gamma_k)$  such that

$$\gamma_k < \gamma_{k-1} < \dots < \gamma_1$$

and

$$\alpha = \omega^{\gamma_1} n_1 + \omega^{\gamma_2} n_2 + \dots + \omega^{\gamma_k} n_k .$$

Proof:

 $\langle 1 \rangle 1$ . For any ordinal  $\alpha$ , there exist a sequence of nonzero natural numbers  $(n_1, \ldots, n_k)$  and sequence of ordinals  $(\gamma_1, \ldots, \gamma_k)$  such that

$$\gamma_k < \gamma_{k-1} < \dots < \gamma_1$$

and

$$\alpha = \omega^{\gamma_1} n_1 + \omega^{\gamma_2} n_2 + \dots + \omega^{\gamma_k} n_k .$$

- $\langle 2 \rangle$ 1. Let:  $\alpha$  be an ordinal
- $\langle 2 \rangle$ 2. Assume: as an induction hypothesis that, for all  $\beta < \alpha$ , the theorem holds.
- $\langle 2 \rangle 3$ . Assume: w.l.o.g.  $\alpha \neq 0$
- $\langle 2 \rangle$ 4. Let:  $\gamma_1$ ,  $n_1$ ,  $\rho_1$  be the unique ordinals such that  $0 \neq n_1 < \omega$ ,  $\rho_1 < \omega^{\gamma_1}$ , and  $\alpha = \omega^{\gamma_1} n_1 + \rho_1$
- $\langle 2 \rangle$ 5. Let:  $(\gamma_2, \ldots, \gamma_k)$  and  $(n_2, \ldots, n_k)$  be sequences such that  $\gamma_k < \gamma_{k-1} < \cdots < \gamma_2$  and  $\rho_1 = \omega^{\gamma_2} n_2 + \cdots + \omega^{\gamma_k} n_k$
- $\langle 2 \rangle 6. \ \gamma_2 < \gamma_1$

PROOF: Since  $\omega^{\gamma_2} \le \rho_1 < \omega^{\gamma_1}$ 

 $\langle 1 \rangle 2$ . If

$$\gamma_k < \gamma_{k-1} < \dots < \gamma_1 \gamma'_k < \gamma'_{k-1} < \dots < \gamma'_1$$

and

$$\omega^{\gamma_1} n_1 + \omega^{\gamma_2} n_2 + \dots + \omega^{\gamma_k} n_k = \omega^{\gamma'_1} n'_1 + \omega^{\gamma'_2} n'_2 + \dots + \omega^{\gamma'_k} n'_k$$
then  $\gamma_i = \gamma'_i$  for all  $i$  and  $n_i = n'_i$  for all  $i$ 

PROOF: Prove by induction on i using the Logarithm Theorem.  $\square$ 

**Definition 9.8.2** (Cantor Normal Form). For any ordinal  $\alpha$ , the *Cantor normal* form of  $\alpha$  is the expression  $\alpha = \omega^{\gamma_1} n_1 + \cdots + \omega^{\gamma_k} n_k$  such that  $n_1, \ldots, n_k$  are nonzero natural numbers and  $\gamma_k < \gamma_{k-1} < \cdots < \gamma_1$ .

# Chapter 10

# The Cumulative Hierarchy

**Definition 10.0.1** (Cumulative Hierarchy). Define the function  $V: \mathbf{On} \to \mathbf{V}$  by transfinite recursion thus:

$$V_{\alpha} = \bigcup_{\beta < \alpha} \mathcal{P}V_{\beta}$$

**Proposition 10.0.2.** For all  $\alpha \in \mathbf{On}$ ,  $V_{\alpha}$  is a transitive set.

Proof:

 $\langle 1 \rangle 1$ . Let:  $\alpha \in \mathbf{On}$ 

 $\langle 1 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall \beta < \alpha. V_{\beta}$  is a transitive set.

 $\langle 1 \rangle 3$ . For all  $\beta < \alpha$ ,  $\mathcal{P}V_{\beta}$  is a transitive set.

PROOF: Proposition 1.6.4.  $\langle 1 \rangle 4$ .  $V_{\alpha}$  is a transitive set. PROOF: Proposition 1.6.3.

| Proof: Propositio

**Proposition 10.0.3.** For any ordinals  $\alpha$  and  $\beta$ , if  $\beta < \alpha$  then  $V_{\beta} \subseteq V_{\alpha}$ .

Proof: Since  $V_{\beta} \in \mathcal{P}V_{\beta} \subseteq V_{\alpha}$  and  $V_{\alpha}$  is a transitive set.  $\square$ 

Theorem 10.0.4.

1. 
$$V_0 = \emptyset$$

2. 
$$\forall \alpha \in \mathbf{On}.V_{\alpha^+} = \mathcal{P}V_{\alpha}$$

3. For any limit ordinal  $\lambda$ ,  $V_{\lambda} = \bigcup_{\alpha < \lambda} V_{\alpha}$ .

Proof:

 $\langle 1 \rangle 1. \ V_0 = \emptyset$ 

PROOF: Immediate from definition.

 $\langle 1 \rangle 2. \ \forall \alpha \in \mathbf{On}.V_{\alpha^+} = \mathcal{P}V_{\alpha}$ 

Proof:

- $\langle 2 \rangle 1$ . Let:  $\alpha \in \mathbf{On}$
- $\langle 2 \rangle 2$ . For all  $\beta < \alpha$  we have  $\mathcal{P}V_{\beta} \subseteq \mathcal{P}V_{\alpha}$ PROOF: Propositions 1.5.8 and 10.0.3.
- $\langle 2 \rangle 3. \ V_{\alpha^+} = \mathcal{P} V_{\alpha}$

$$V_{\alpha^{+}} = \bigcup_{eta < \alpha^{+}} \mathcal{P}V_{eta}$$

$$= \bigcup_{eta < \alpha} \mathcal{P}V_{eta} \cup \mathcal{P}V_{\alpha}$$
 $\mathcal{P}V_{\alpha}$ 

 $\langle 1 \rangle$ 3. For any limit ordinal  $\lambda$ ,  $V_{\lambda} = \bigcup_{\alpha < \lambda} V_{\alpha}$ 

Proof:

 $\langle 2 \rangle 1. \ V_{\lambda} \subseteq \bigcup_{\alpha < \lambda} V_{\alpha}$ 

Proof:

$$V_{\lambda} = \bigcup_{\alpha < \lambda} \mathcal{P}V_{\alpha}$$

$$= \bigcup_{\alpha < \lambda} V_{\alpha^{+}} \qquad (\langle 1 \rangle 2)$$

$$\subseteq \bigcup_{\alpha < \lambda} V_{\alpha}$$

 $\langle 2 \rangle 2. \bigcup_{\alpha < \lambda} V_{\alpha} \subseteq V_{\lambda}$ PROOF: Proposition 10.0.3.

**Proposition 10.0.5.** For every set A, there exists an ordinal  $\alpha$  such that  $A \in$  $V_{\alpha}$ .

## Proof:

- $\langle 1 \rangle 1$ . Let us say a set A is grounded iff there exists an ordinal  $\alpha$  such that  $A \in V_{\alpha}$ .
- $\langle 1 \rangle 2$ . For any set A, if every element of A is grounded, then A is grounded.
  - $\langle 2 \rangle 1$ . Let: A be a set.
  - $\langle 2 \rangle 2$ .  $S = \{ \alpha \mid \exists a \in A.\alpha \text{ is the least ordinal such that } a \in V_{\alpha} \}$

PROOF: S is a set by an Axiom of Replacement.

- $\langle 2 \rangle 3$ . Let:  $\beta = \sup S$
- $\langle 2 \rangle 4$ .  $A \subseteq V_{\beta}$ 
  - $\langle 3 \rangle 1$ . Let:  $a \in A$
  - $\langle 3 \rangle 2$ . Let:  $\alpha$  be the least ordinal such that  $a \in V_{\beta}$
  - $\langle 3 \rangle 3. \ \alpha \in S$
  - $\langle 3 \rangle 4. \ \alpha \leq \beta$
  - $\langle 3 \rangle 5. \ a \in V_{\beta}$
- $\langle 2 \rangle 5. \ A \in V_{\beta^+}$
- $\langle 1 \rangle 3$ . Assume: for a contradiction there exists an ungrounded set.
- $\langle 1 \rangle 4$ . PICK a transitive set B that has an ungrounded member.

PROOF: Pick a transitive set c, and take B to be the transitive closure of  $\{c\}$ .

```
\begin{array}{l} \langle 1 \rangle \text{5. Let: } A = \{x \in B \mid x \text{ is ungrounded}\} \\ \langle 1 \rangle \text{6. Pick } m \in A \text{ such that } m \cap A = \emptyset \\ \text{Proof: Axiom of Regularity.} \\ \langle 1 \rangle \text{7. Every member of } m \text{ is grounded.} \\ \langle 2 \rangle \text{1. Assume: for a contradiction } x \in m \text{ is ungrounded.} \\ \langle 2 \rangle \text{2. } x \in B \\ \text{Proof: Since } B \text{ is transitive } (\langle 1 \rangle 4). \\ \langle 2 \rangle \text{3. } x \in A \\ \text{Proof: } \langle 1 \rangle \text{5} \\ \langle 2 \rangle \text{4. Q.E.D.} \\ \text{Proof: This contradicts } \langle 1 \rangle \text{6.} \\ \langle 1 \rangle \text{8. } m \text{ is grounded.} \\ \text{Proof: } \langle 1 \rangle \text{2} \\ \langle 1 \rangle \text{9. Q.E.D.} \\ \text{Proof: This contradicts } \langle 1 \rangle \text{6.} \\ \end{array}
```

**Definition 10.0.6** (Rank). The *rank* of a set A is the least ordinal  $\alpha$  such that  $A \in V_{\alpha^+}$ .

**Proposition 10.0.7.** For any set A we have

$$\operatorname{rank} A = \bigcup_{a \in A} (\operatorname{rank} a)^+$$

```
PROOF:  \langle 1 \rangle 1. \text{ Let: } \alpha = \bigcup_{a \in A} (\operatorname{rank} a)^+ \\ \langle 1 \rangle 2. A \subseteq V_{\alpha} \\ \langle 2 \rangle 1. \text{ Let: } a \in A \\ \langle 2 \rangle 2. a \in V_{(\operatorname{rank} a)^+} \\ \langle 2 \rangle 3. a \in V_{\alpha} \\ \langle 1 \rangle 3. A \in V_{\alpha^+} \\ \langle 1 \rangle 4. \text{ If } A \subseteq V_{\beta} \text{ then } \alpha \leq \beta \\ \langle 2 \rangle 1. \text{ Assume: } A \subseteq V_{\beta} \\ \langle 2 \rangle 2. \text{ For all } a \in A \text{ we have } (\operatorname{rank} a)^+ \leq \beta \\ \text{PROOF: Since } a \in V_{\beta}. \\ \langle 2 \rangle 3. \alpha \leq \beta
```

**Corollary 10.0.7.1.** *For any sets a and b, if*  $a \in b$  *then* rank  $a < \operatorname{rank} b$ .

**Proposition 10.0.8.** For any ordinal number  $\alpha$  we have rank  $\alpha = \alpha$ .

#### Proof:

- $\langle 1 \rangle 1$ . Let:  $\alpha$  be an ordinal.
- $\langle 1 \rangle 2$ . Assume: as transfinite induction hypothesis  $\forall \beta < \alpha$ . rank  $\beta = \beta$
- $\langle 1 \rangle 3$ . rank  $\alpha = \bigcup_{\beta < \alpha} \beta^+$

Proof:

$$\operatorname{rank} \alpha = \bigcup_{\beta < \alpha} (\operatorname{rank} \beta)^+$$
$$= \bigcup_{\beta < \alpha} \beta^+$$

 $\langle 1 \rangle 4$ .  $\bigcup_{\beta < \alpha} \beta^+ \le \alpha$ PROOF: Since for all  $\beta < \alpha$  we have  $\beta^+ \le \alpha$ .

- (1)5.  $\alpha \leq \bigcup_{\beta < \alpha} \beta^+$ (2)1. Let:  $\gamma = \bigcup_{\beta < \alpha} \beta^+$ (2)2. Assume: for a contradiction  $\gamma < \alpha$ (2)3.  $\gamma^+ \leq \bigcup_{\beta < \alpha} \beta^+ = \gamma$ (2)4. Q.E.D.

PROOF: This is a contradiction.

Definition 10.0.9 (Hereditarily Finite). A set is hereditarily finite iff it is in

# Chapter 11

# Models of Set Theory

**Definition 11.0.1** (Relativization). Let  $\sigma$  be a sentence in the language of set theory and  $\mathbf{M}$  a class. The *relativization* of  $\sigma$  to  $\mathbf{M}$  is the sentence  $\sigma^{\mathbf{M}}$  formed by replacing every quantifier  $\forall x$  with  $\forall x \in \mathbf{M}$ , and  $\exists x$  with  $\exists x \in \mathbf{M}$ .

We write 'M is a model of  $\sigma$ ' for the sentence  $\sigma^{\mathbf{M}}$ .

**Theorem Schema 11.0.2.** For any class M, the following is a theorem: If M is a transitive class, then M is a model of the Axiom of Extensionality.

```
Proof:
```

```
\begin{array}{ll} \langle 1 \rangle 1. \text{ Assume: } \mathbf{M} \text{ is a transitive class.} \\ \text{Prove: } \forall x,y \in \mathbf{M} (\forall z \in \mathbf{M} (z \in x \Leftrightarrow z \in y) \Rightarrow x = y) \\ \langle 1 \rangle 2. \text{ Let: } x,y \in \mathbf{M} \\ \langle 1 \rangle 3. \text{ Assume: } \forall z \in \mathbf{M} (z \in x \Leftrightarrow z \in y) \\ \langle 1 \rangle 4. \ \forall z (z \in x \Leftrightarrow z \in y) \\ \text{Proof: Since } z \in x \Rightarrow z \in \mathbf{M} \text{ and } z \in y \Rightarrow z \in \mathbf{M} \text{ by } \langle 1 \rangle 1. \\ \langle 1 \rangle 5. \ x = y \\ \square \end{array}
```

**Theorem 11.0.3.** If  $\alpha$  is a non-zero ordinal then  $V_{\alpha}$  is a model of the statement: The empty class is a set.

### Proof:

```
\begin{array}{ll} \langle 1 \rangle 1. & \text{Let: } \alpha \neq 0 \\ & \text{Prove: } \exists x \in V_{\alpha}. \forall y \in V_{\alpha}. y \notin x \\ \langle 1 \rangle 2. & \emptyset \in V_{\alpha} \\ \langle 1 \rangle 3. & \forall y \in V_{\alpha}. y \notin \emptyset \\ & \Box \end{array}
```

**Theorem 11.0.4.** For any limit ordinal  $\lambda$ , we have  $V_{\lambda}$  is a model of the statement: for any sets a and b, the class  $\{a,b\}$  is a set.

#### Proof:

 $\langle 1 \rangle 1$ . Let:  $\lambda$  be a limit ordinal.

```
PROVE: \forall a,b \in V_{\lambda}. \exists c \in V_{\lambda}. \forall x \in V_{\lambda} (x \in c \Leftrightarrow x = a \lor x = b) \langle 1 \rangle 2. Let: a,b \in V_{\lambda} \langle 1 \rangle 3. Pick \alpha,\beta < \lambda such that a \in V_{\alpha} and b \in V_{\beta} \langle 1 \rangle 4. Assume: w.l.o.g. \alpha \leq \beta \langle 1 \rangle 5. a,b \in V_{\beta} \langle 1 \rangle 6. \{a,b\} \in V_{\beta+1} \langle 1 \rangle 7. \{a,b\} \in V_{\lambda} \langle 1 \rangle 8. \forall x \in V_{\lambda} (x \in \{a,b\} \Leftrightarrow x = a \lor x = b)
```

**Theorem 11.0.5.** For any ordinal  $\alpha$ , we have  $V_{\alpha}$  is a model of the Union Axiom.

#### Proof:

```
TROOF.  \langle 1 \rangle 1. \text{ Let: } \alpha \text{ be an ordinal.}  Prove:  \forall a \in V_{\alpha}. \exists b \in V_{\alpha}. \forall x \in V_{\alpha} (x \in b \Leftrightarrow \exists y \in V_{\alpha} (x \in y \land y \in a))   \langle 1 \rangle 2. \text{ Let: } a \in V_{\alpha}   \langle 1 \rangle 3. \text{ Pick } \beta < \alpha \text{ such that } a \subseteq V_{\beta}   \langle 1 \rangle 4. \bigcup a \subseteq V_{\beta}  Proof:  V_{\beta} \text{ is a transitive set.}   \langle 1 \rangle 5. \bigcup a \in V_{\alpha}   \langle 1 \rangle 6. \forall x \in V_{\alpha} (x \in \bigcup a \Leftrightarrow \exists y \in V_{\alpha} (x \in y \land y \in a))  Proof:  V_{\alpha} \text{ is a transitive set.}
```

**Theorem 11.0.6.** For any limit ordinal  $\lambda$ , we have  $V_{\lambda}$  is a model of the Power Set Axiom.

#### Proof:

**Theorem Schema 11.0.7.** For any property  $P[x, y_1, ..., y_n]$ , the following is a theorem:

For any ordinal  $\alpha$ , the set  $V_{\alpha}$  is a model of the statement: for any sets  $a_1$ , ...,  $a_n$ , B, the class  $\{x \in B \mid P[x, a_1, ..., a_n]\}$  is a set.

## Proof:

 $\langle 1 \rangle 1$ . Let:  $\alpha$  be an ordinal.  $\langle 1 \rangle 2$ . Let:  $a_1, \ldots, a_n, B \in V_{\alpha}$   $\langle 1 \rangle 3$ . Let:  $C = \{x \in B \mid P[x, a_1, \ldots, a_n]^{V_{\alpha}}\}$   $\langle 1 \rangle 4$ .  $C \in V_{\alpha}$ 

```
\langle 1 \rangle 5. \ \forall x \in V_{\alpha}(x \in C \Leftrightarrow x \in B \land P[x, a_1, \dots, a_n]^{V_{\alpha}})
```

**Theorem 11.0.8.** For any ordinal  $\alpha > \omega$ , we have:  $V_{\alpha}$  is a model of the Axiom of Infinity.

### Proof:

- $\langle 1 \rangle 1$ . Let:  $\alpha > \omega$
- $\langle 1 \rangle 2. \ \mathbb{N} \in V_{\alpha}$
- $\langle 1 \rangle 3. \ \exists e \in V_{\alpha} (e \in \mathbb{N} \land \forall x \in V_{\alpha}.x \notin e)$
- $\langle 1 \rangle 4. \ \forall x \in V_{\alpha}(x \in \mathbb{N} \Rightarrow \exists y \in V_{\alpha} \forall z \in V_{\alpha}(z \in y \Leftrightarrow z \in x \lor z = x))$

**Theorem 11.0.9.** For any ordinal  $\alpha$ , we have  $V_{\alpha}$  is a model of the Axiom of Choice.

### Proof:

- $\langle 1 \rangle 1$ . Let:  $\alpha$  be an ordinal.
- $\langle 1 \rangle 2$ . Let:  $A \in V_{\alpha}$
- $\langle 1 \rangle 3$ . Assume:  $\forall x \in V_{\alpha} (x \in A \Rightarrow \exists y \in V_{\alpha}. y \in A)$
- $\langle 1 \rangle 4$ . Assume:  $\forall x, y, z \in V_{\alpha} (x \in A \land y \in A \land z \in x \land z \in y \Rightarrow x = y)$
- $\langle 1 \rangle 5$ . A is a set of pairwise disjoint nonempty sets.
- $\langle 1 \rangle 6$ . Pick c such that, for all  $x \in A$ ,  $x \cap c = \emptyset$
- $\langle 1 \rangle 7. \ c \cap \bigcup A \in V_{\alpha}$
- $(1) 8. \ \forall x \in V_{\alpha}(x \in A \Rightarrow \exists y \in V_{\alpha} \forall z \in V_{\alpha}(z \in x \land z \in c \cap \bigcup A \Leftrightarrow z = y))$

**Theorem 11.0.10.** For any ordinal  $\alpha$ , we have  $V_{\alpha}$  is a model of the Axiom of Regularity.

### Proof:

- $\langle 1 \rangle 1$ . Let:  $\alpha$  be an ordinal.
- $\langle 1 \rangle 2$ . Let:  $A \in V_{\alpha}$
- $\langle 1 \rangle 3$ . Assume:  $\exists x \in V_{\alpha}.x \in A$
- $\langle 1 \rangle 4$ . Pick  $m \in A$  of least rank.
- $\langle 1 \rangle 5. \ m \in V_{\alpha}$
- $\langle 1 \rangle 6. \ \neg \exists x \in V_{\alpha} (x \in m \land x \in A)$

**Theorem Schema 11.0.11.** For any axiom  $\alpha$  of Zermelo set theory, the following is a theorem:

For any limit ordinal  $\lambda > \omega$ , we have  $V_{\lambda}$  is a model of  $\alpha$ .

PROOF: Theorems 11.0.2, 11.0.3, 11.0.4, 11.0.5, 11.0.6, 11.0.7, 11.0.8, 11.0.9, 11.0.10.  $\Box$ 

**Corollary Schema 11.0.11.1.** for any axiom  $\alpha$  of Zermelo set theory, the following is a theorem:

 $V_{\omega 2}$  is a model of  $\alpha$ .

**Lemma 11.0.12.** There exists a well-ordered structure in  $V_{\omega 2}$  whose ordinal is not in  $V_{\omega 2}$ .

PROOF: Take the well-ordered set  $\mathbb{N} \times \{0,1\}$  whose ordinal is  $\omega 2$ .  $\square$ 

Corollary Schema 11.0.12.1. There exists an instance  $\alpha$  of the Axiom Schema of Replacement such that the following is a theorem:

 $V_{\omega 2}$  is not a model of  $\alpha$ .

## Chapter 12

## Infinite Cardinals

### 12.1 Arithmetic of Infinite Cardinals

**Proposition 12.1.1.** For any infinite cardinal  $\kappa$  we have  $\kappa \kappa = \kappa$ .

```
Proof:
\langle 1 \rangle 1. PICK a set B with |B| = \kappa
\langle 1 \rangle 2. Let: \mathcal{H} = \{ f \mid f = \emptyset \lor \exists A \subseteq B. (A \text{ is infinite} \land f : A \times A \approx A \}
\langle 1 \rangle 3. For any chain \mathcal{C} \subseteq \mathcal{H} we have \bigcup \mathcal{C} \in \mathcal{H}
    \langle 2 \rangle 1. Let: \mathcal{C} \subseteq \mathcal{H} be a chain.
    \langle 2 \rangle 2. Assume: w.l.o.g. \mathcal C has a nonempty element.
    \langle 2 \rangle 3. \bigcup \mathcal{C} is a function.
         \langle 3 \rangle 1. Assume: (x,y), (x,z) \in \bigcup \mathcal{C}
         \langle 3 \rangle 2. PICK f, g \in \mathcal{C} such that f(x) = y and g(x) = z
         \langle 3 \rangle 3. Assume: w.l.o.g. f \subseteq g
         \langle 3 \rangle 4. \ y=z
    \langle 2 \rangle 4. \bigcup \mathcal{C} is injective.
         Proof: Similar.
     \langle 2 \rangle5. Let: A = \operatorname{ran} \bigcup \mathcal{C}
    \langle 2 \rangle 6. A is infinite.
         \langle 3 \rangle 1. PICK a nonzero f \in \mathcal{C}
         \langle 3 \rangle 2. Let: A' be the infinite subset of B such that f: A'^2 \approx A'
         \langle 3 \rangle 3. \ A' \subseteq A
    \langle 2 \rangle 7. dom \bigcup \mathcal{C} = A^2
         \langle 3 \rangle 1. Let: x, y \in A
         \langle 3 \rangle 2. PICK f, g \in \mathcal{C} such that x \in \operatorname{ran} f and y \in \operatorname{ran} g
         \langle 3 \rangle 3. Assume: w.l.o.g. f \subseteq g
         \langle 3 \rangle 4. Let: A' be the infinite subset of B such that g:A'^2 \approx A'
         \langle 3 \rangle 5. \ x, y \in A'
         \langle 3 \rangle 6. \ (x,y) \in \text{dom } g
         \langle 3 \rangle 7. \ (x,y) \in \operatorname{dom} \bigcup \mathcal{C}
    \langle 2 \rangle 8. \bigcup \mathcal{C} \in \mathcal{H}
```

- $\langle 1 \rangle 4$ . Pick a maximal  $f_0 \in \mathcal{H}$
- $\langle 1 \rangle 5. \ f_0 \neq \emptyset$ 
  - $\langle 2 \rangle 1$ . PICK a countably infinite subset A of B.

Proof: Proposition 9.2.10.

 $\langle 2 \rangle 2$ . Pick a bijection  $f: A^2 \approx A$ 

Proof: Proposition 9.2.9.

- $\langle 2 \rangle 3. \ \emptyset \subseteq f \in \mathcal{H}$
- $\langle 2 \rangle 4$ .  $\emptyset$  is not maximal in  $\mathcal{H}$
- $\langle 1 \rangle 6$ . Let:  $A_0$  be the infinite subset of B such that  $f_0: A_0^2 \approx A_0$
- $\langle 1 \rangle 7$ . Let:  $\lambda = |A_0|$
- $\langle 1 \rangle 8$ .  $\lambda$  is infinite.
- $\langle 1 \rangle 9. \ \lambda^2 = \lambda$
- $\langle 1 \rangle 10. \ \lambda = \kappa$ 
  - $\langle 2 \rangle 1$ . Assume: for a contradiction  $\lambda < \kappa$
  - $\langle 2 \rangle 2$ .  $\lambda \leq |B A_0|$
  - $\langle 2 \rangle 3$ . Pick a subset  $D \subseteq B A_0$  with  $|D| = \lambda$
  - $\langle 2 \rangle 4$ .  $(A_0 \cup D)^2 = A_0^2 \cup (A_0 \times D) \cup (D \times A_0) \cup D^2$  $\langle 2 \rangle 5$ . Let:  $C = (A_0 \times D) \cup (D \times A_0) \cup D^2$

  - $\langle 2 \rangle 6. \ |C| = \lambda$

Proof:

$$|(A_0 \times D) \cup (D \times A_0) \cup D^2| = \lambda^2 + \lambda^2 + \lambda^2$$

$$= \lambda + \lambda + \lambda \qquad (\langle 1 \rangle 9)$$

$$= 3\lambda$$

$$\leq \lambda \cdot \lambda$$

$$= \lambda \qquad (\langle 1 \rangle 9)$$

- $\langle 2 \rangle$ 7. Pick a bijection  $g: C \approx D$
- $\langle 2 \rangle 8.$   $f_0 \cup g : (A_0 \cup D)^2 \approx A_0 \cup D$
- $\langle 2 \rangle 9$ . Q.E.D.

PROOF: This contradicts the maximality of  $f_0$ .

**Theorem 12.1.2** (Absorpution Law of Cardinal Arithmetic). Let  $\kappa$  and  $\lambda$  be nonzero cardinal numbers such that at least one is infinite. Then

$$\kappa + \lambda = \kappa \lambda = \max(\kappa, \lambda)$$

- $\langle 1 \rangle 1$ . Assume: w.l.o.g.  $\lambda \leq \kappa$
- $\langle 1 \rangle 2$ .  $\kappa + \lambda = \kappa \lambda = \kappa$

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Proof:

$$\begin{split} \kappa &\leq \kappa + \lambda \\ &\leq \kappa + \kappa \\ &= 2\kappa \\ &\leq \kappa \lambda \\ &\leq \kappa \kappa \\ &= \kappa \end{split} \tag{Proposition 12.1.1}$$

## 12.2 Alephs

**Definition 12.2.1** (Aleph). Let  $\aleph$  be the unique order isomorphism between **On** and the class of infinite cardinals.

**Proposition 12.2.2.** The operation  $\aleph$  is normal.

Proof: Proposition 7.6.8 and Lemma 8.2.7.  $\square$ 

**Definition 12.2.3** (Continuum Hypothesis). The *continuum hypothesis* is the statement that  $\aleph_1 = 2^{\aleph_0}$ .

**Definition 12.2.4** (Generalised Continuum Hypothesis). The generalised continuum hypothesis is the statement that, for all  $\alpha$ ,  $\aleph_{\alpha^+} = 2^{\aleph_{\alpha}}$ .

### 12.3 Beths

**Definition 12.3.1** (Beth). Define the operation  $\beth: \mathbf{On} \to \mathbf{Card}$  by transfinite recursion as follows:

$$\begin{split} & \beth_0 := \aleph_0 \\ & \beth_{\alpha^+} := 2^{\beth_\alpha} \\ & \beth_\lambda := \bigcup_{\alpha < \lambda} \beth_\alpha \end{split} \qquad (\lambda \text{ a limit ordinal})$$

**Proposition 12.3.2.**  $\supset$  *is a normal operation.* 

PROOF: It is continuous by definition, and  $\beth_{\alpha} < \beth_{\alpha^+}$  by Cantor's Theorem.  $\square$ 

**Proposition 12.3.3.** The continuum hypothesis is equivalent to the statement  $\beth_1 = \aleph_1$ .

The generalised continuum hypothesis is equivalent to the statement  $\beth = \alpha$ .

Proof: Immediate from definitions.  $\square$ 

**Lemma 12.3.4.** For any ordinal number  $\alpha$ , we have  $|V_{\omega+\alpha}| = \beth_{\alpha}$ .

 $\langle 1 \rangle 1. |V_{\omega}| = \beth_0$ 

PROOF: Since  $V_{\omega}$  is the union of  $\aleph_0$  finite sets of increasing size.

 $\langle 1 \rangle$ 2. For any ordinal  $\alpha$ , if  $|V_{\omega+\alpha}| = \beth_{\alpha}$  then  $|V_{\omega+\alpha+1}| = \beth_{\alpha+1}$  PROOF: Since  $V_{\omega+\alpha+1} = \mathcal{P}V_{\omega+\alpha}$ .

 $\langle 1 \rangle 3$ . For any limit ordinal  $\lambda$ , if  $\forall \alpha < \lambda . |V_{\omega+\alpha}| = \beth_{\alpha}$  then  $|V_{\omega+\lambda}| = \beth_{\lambda}$ . Proof:

$$|V_{\omega+\lambda}| = \left| \bigcup_{\alpha < \lambda} V_{\omega+\alpha} \right|$$

$$= \sup_{\alpha < \lambda} |V_{\omega+\alpha}|$$

$$= \sup_{\alpha < \lambda} \beth_{\alpha}$$

$$= \beth_{\lambda}$$

## 12.4 Cofinality

**Definition 12.4.1** (Cofinal). Let  $\lambda$  be a limit ordinal and S a set of ordinals smaller than  $\lambda$ . Then S is *cofinal* in  $\lambda$  if and only if  $\lambda = \sup S$ .

**Definition 12.4.2** (Cofinality). For any ordinal  $\alpha$ , define the *cofinality* of  $\alpha$ , of  $\alpha$ , as follows:

- cf 0 = 0
- For any ordinal  $\alpha$ , cf  $\alpha^+ = 1$
- For any limit ordinal  $\lambda$ , cf  $\lambda$  is the smallest cardinal such that there exists a set S of ordinals cofinal in  $\lambda$  with  $|S| = \operatorname{cf} \lambda$ .

**Definition 12.4.3** (Regular). A cardinal  $\kappa$  is regular iff cf  $\kappa = \kappa$ ; otherwise it is singular.

**Proposition 12.4.4.**  $\aleph_0$  is regular.

PROOF:  $\aleph_0$  is not the supremum of  $< \aleph_0$  smaller ordinals, because a finite union of finite ordinals is finite.  $\square$ 

**Proposition 12.4.5.** For every ordinal  $\alpha$ ,  $\aleph_{\alpha+1}$  is regular.

PROOF: If S is a set of ordinals with  $|S| < \aleph_{\alpha+1}$  and  $\forall \beta \in S.\beta < \aleph_{\alpha+1}$ , then we have  $|S| \leq \aleph_{\alpha}$  and  $\forall \beta \in S.\beta \leq \aleph_{\alpha}$ , hence

$$\left| \bigcup S \right| \leq \aleph_{\alpha}^{2} \qquad \qquad \text{(Proposition 8.2.6)}$$
  
=  $\aleph_{\alpha}$  \quad \text{(Proposition 12.1.1)}

**Proposition Schema 12.4.6.** For any class  $\mathbf{T}$ , the following is a theorem. Assume  $\mathbf{T}: \mathbf{On} \to \mathbf{On}$  is a normal operation. For any limit ordinal  $\lambda$  we have  $\operatorname{cf} \mathbf{T}(\lambda) = \operatorname{cf} \lambda$ .

```
Proof:
\langle 1 \rangle 1. cf \mathbf{T}(\lambda) \leq \operatorname{cf} \lambda
     \langle 2 \rangle 1. Pick a set S of ordinals \langle \lambda \rangle with |S| = \operatorname{cf} \lambda and \sup S = \lambda
     \langle 2 \rangle 2. \mathbf{T}(\lambda) = \sup_{\alpha \in S} \mathbf{T}(\alpha)
          PROOF: Theorem 7.6.7.
\langle 1 \rangle 2. cf \lambda < cf \mathbf{T}(\lambda)
     \langle 2 \rangle 1. Pick a set A of ordinals \langle \mathbf{T}(\lambda) \rangle such that |A| = \operatorname{cf} \mathbf{T}(\lambda) and \sup A = \operatorname{cf} \mathbf{T}(\lambda)
                   \mathbf{T}(\lambda)
     \langle 2 \rangle 2. Let: B = \{ \gamma < \lambda \mid \exists \alpha \in A. |\alpha| = \mathbf{T}(\gamma) \}
     \langle 2 \rangle 3. |B| \leq |A| = \operatorname{cf} \mathbf{T}(\lambda)
                  Prove: \sup B = \lambda
     \langle 2 \rangle 4. \ \forall \alpha \in A. |\alpha| \leq \mathbf{T}(\sup B)
     \langle 2 \rangle 5. \ \forall \alpha \in A.\alpha < \mathbf{T}(\sup B + 1)
     \langle 2 \rangle 6. \aleph_{\lambda} = \sup A \leq \mathbf{T}(\sup B + 1)
     \langle 2 \rangle 7. \lambda \leq \sup B + 1
     \langle 2 \rangle 8. \ \lambda \leq \sup B
          PROOF: \lambda is a limit ordinal.
      \langle 2 \rangle 9. sup B = \lambda
П
```

Corollary 12.4.6.1.  $\aleph_{\omega}$  is singular.

PROOF:  $\operatorname{cf} \aleph_{\omega} = \operatorname{cf} \aleph_0 = \aleph_0$ .  $\square$ 

Corollary 12.4.6.2. The operation of is not strictly monotone or continuous.

PROOF: cf  $\aleph_{\omega}$  < cf  $\aleph_1$ 

**Definition 12.4.7** (Weakly Inaccessible). A cardinal is *weakly inaccessible* iff it is  $\aleph_{\lambda}$  for some limit ordinal  $\lambda$  and regular.

**Lemma 12.4.8.** Let  $\lambda$  be a limit ordinal. Then there exists a strictly increasing of  $\lambda$ -sequence that converges to  $\lambda$ .

### Proof:

- $\langle 1 \rangle 1$ . Pick a set S of ordinals  $\langle \lambda \rangle$  with  $|S| = \operatorname{cf} \lambda$  and  $\sup S = \lambda$
- $\langle 1 \rangle 2$ . Pick a bijection  $a : \text{cf } \lambda \approx S$
- $\langle 1 \rangle 3$ . PICK a strictly increasing subsequence  $(b_{\delta})_{\delta < \beta}$  of a that converges to  $\lambda$ . PROOF: Lemma 7.8.5.

 $\langle 1 \rangle 4$ .  $\beta = \operatorname{cf} \lambda$ 

PROOF: By minimiality of cf  $\lambda$ .

Corollary 12.4.8.1. Let  $\lambda$  be a limit ordinal. Then cf  $\lambda$  is the least ordinal such that there exists a strictly increasing cf  $\lambda$ -sequence that converges to  $\lambda$ .

**Proposition 12.4.9.** For any ordinal  $\lambda$ , cf  $\lambda$  is a regular cardinal.

```
\langle 1 \rangle 1. Let: \lambda be an ordinal.
```

 $\langle 1 \rangle 2$ . Assume: w.l.o.g.  $\lambda$  is a limit ordinal.

 $\langle 1 \rangle 3$ . Pick a strictly increasing sequence  $(a_{\alpha})_{\alpha < \text{cf } \lambda}$  that converges to  $\lambda$ .

(1)4. Let: S be a set of ordinals  $\langle \operatorname{cf} \lambda \operatorname{such that} | S | = \operatorname{cf} \operatorname{cf} \lambda \operatorname{and sup} S = \operatorname{cf} \lambda$ .

 $\langle 1 \rangle 5$ . Let:  $a(S) = \{a_{\alpha} \mid \alpha \in S\}$ 

 $\langle 1 \rangle 6$ . a(S) is cofinal in  $\lambda$ .

 $\langle 2 \rangle 1$ . Let:  $\beta < \lambda$ 

 $\langle 2 \rangle 2$ . Pick  $\gamma < \text{cf } \lambda \text{ such that } \beta < a_{\gamma}$ 

 $\langle 2 \rangle 3$ . Pick  $\delta \in S$  such that  $\gamma < \delta$ 

 $\langle 2 \rangle 4$ .  $a_{\delta} \in a(S)$  and  $\beta < a_{\gamma} < a_{\delta}$ 

 $\langle 1 \rangle 7$ . cf  $\lambda \leq$  cf cf  $\lambda$ 

PROOF: Since a(S) is a set of ordinals  $<\lambda$  with  $|a(S)|=\mathrm{cf}\,\mathrm{cf}\,\lambda$  and  $\sup a(S)=\lambda$ .

 $\langle 1 \rangle 8$ . cf cf  $\lambda = \text{cf } \lambda$ 

**Theorem 12.4.10.** Let  $\lambda$  be an infinite cardinal. Then cf  $\lambda$  is the least cardinal such that  $\lambda$  can be partitioned into cf  $\lambda$  sets, each of cardinality  $< \lambda$ .

#### Proof

- $\langle 1 \rangle 1$ .  $\lambda$  can be partitioned into cf  $\lambda$  sets, each of cardinality  $\langle \lambda \rangle$ 
  - $\langle 2 \rangle$ 1. PICK a strictly increasing sequence of ordinlas  $(a_{\alpha})_{\alpha < \operatorname{cf} \lambda}$  that converges to  $\lambda$
  - $\langle 2 \rangle 2$ .  $\{ \{ \beta \mid a_{\alpha} \leq \beta < a_{\alpha+1} \} \mid \alpha < \text{cf } \lambda \} \text{ is a partition of } \lambda \text{ into cf } \lambda \text{ sets, each of cardinality } < \lambda$
- $\langle 1 \rangle 2$ . If  $\lambda$  can be partitioned into  $\kappa$  sets, each of cardinality  $\langle \lambda$ , then cf  $\lambda \leq \kappa$ .
  - $\langle 2 \rangle 1$ . Let:  $\mathcal{A}$  be a partition of  $\lambda$  into sets of cardinality  $\langle \lambda \rangle$
  - $\langle 2 \rangle 2$ . Let:  $\kappa = |P|$
  - $\langle 2 \rangle 3$ . Pick a bijection  $A : \kappa \approx P$
  - $\langle 2 \rangle 4. \ \lambda = \bigcup_{\xi < \kappa} A(\xi)$
  - $\langle 2 \rangle 5$ . For all  $\xi < \kappa$  we have  $|A(\xi)| < \lambda$
  - $\langle 2 \rangle 6$ . Let:  $\mu = \sup_{\xi < \kappa} |A(\xi)|$
  - $\langle 2 \rangle 7. \ \mu \leq \lambda$
  - $\langle 2 \rangle 8$ . For all  $\xi < \kappa$  we have  $|A(\xi)| \leq \mu$
  - $\langle 2 \rangle 9. \ \lambda \leq \mu \kappa$

Proof: Proposition 8.2.6.

 $\langle 2 \rangle 10$ . Assume: w.l.o.g.  $\kappa < \lambda$ 

PROOF: If  $\lambda \leq \kappa$  then cf  $\lambda \leq \kappa$  since cf  $\lambda \leq \lambda$ .

 $\langle 2 \rangle 11. \ \lambda = \mu$ 

Proof:

$$\lambda \leq \mu \kappa \qquad (\langle 2 \rangle 9)$$
  
$$\leq \lambda \lambda \qquad (\langle 2 \rangle 7, \langle 2 \rangle 10)$$
  
$$= \lambda \qquad (Proposition 12.1.1)$$

 $\langle 2 \rangle 12$ .  $\{|A(\xi)| \mid \xi < \kappa\}$  is a set of  $\leq \kappa$  ordinals all  $< \lambda$  whose supremum is  $\lambda \langle 2 \rangle 13$ . cf  $\lambda \leq \kappa$ 

**Theorem 12.4.11** (König). For any infinite cardinal  $\kappa$  we have  $\kappa < \operatorname{cf} 2^{\kappa}$ .

```
Proof:
```

- $\langle 1 \rangle 1$ . Assume: for a contradiction of  $2^{\kappa} \leq \kappa$
- $\langle 1 \rangle 2$ . Let:  $S = 2^{\kappa}$
- $\langle 1 \rangle 3$ . Pick a partition  $\{A_{\xi} \mid \xi < \kappa\}$  of  $S^{\kappa}$  with  $\forall \xi < \kappa . |A_{\xi}| < 2^{\kappa}$ .

PROOF: Theorem 12.4.10.

 $\langle 1 \rangle 4. \ \forall \xi < \kappa. \{ g(\xi) \mid g \in A_{\xi} \} \subsetneq S$ 

PROOF: We do not have equality because  $|\{g(\xi) \mid g \in A_{\xi}\}| \leq |A_{\xi}| < 2^{\kappa}$ .

- $\langle 1 \rangle 5$ . For all  $\xi < \kappa$ , choose  $s_{\xi} \in S \{g(\xi) \mid g \in A_{\xi}\}$
- $\langle 1 \rangle 6. \ s \in S^{\kappa}$
- $\langle 1 \rangle 7$ . For all  $\xi < \kappa$  we have  $s \notin A_{\xi}$

PROOF: Since for all  $\xi < \kappa$  and  $g \in A_{\xi}$  we have  $s_{\xi}(\xi) \neq g(\xi)$ .

 $\langle 1 \rangle 8$ . Q.E.D.

PROOF: This contradicts  $\langle 1 \rangle 3$ .

### Corollary 12.4.11.1.

$$2^{\aleph_0} \neq \aleph_\omega$$

**Proposition 12.4.12.** For any ordinal  $\alpha$ , we have cf  $\alpha$  is the least cardinal such that  $\alpha$  is the strict supremum of cf  $\alpha$  smaller ordinals.

#### Proof:

```
\langle 1 \rangle 1. Case: \alpha = 0
```

PROOF: Since  $0 = \text{ssup } \emptyset$ .

 $\langle 1 \rangle 2$ . Case:  $\alpha = \beta^+$ 

PROOF: Since  $\beta^+ = \text{ssup}\{\beta\}$ .

- $\langle 1 \rangle 3$ . Case:  $\alpha$  is a limit ordinal.
  - $\langle 2 \rangle 1$ . There exists a set S of ordinals  $\langle \alpha \rangle$  such that  $|S| = cf \alpha$  and  $\alpha = ssup S$ .
    - $\langle 3 \rangle$ 1. PICK a set S of ordinals  $< \alpha$  such that  $|S| = \text{cf } \alpha$  and  $\sup S = \alpha$  PROVE:  $\alpha = \text{ssup } S$
    - $\langle 3 \rangle 2. \ \forall \beta \in S.\beta < \alpha$
    - $\langle 3 \rangle 3$ . For any ordinal  $\gamma$ , if  $\forall \beta \in S.\beta < \gamma$  then  $\alpha \leq \gamma$
  - $\langle 2 \rangle 2$ . If T is a set of ordinals  $\langle \alpha \rangle$  such that  $\alpha = \sup T$ , then cf  $\alpha \leq |T|$ .
    - $\langle 3 \rangle 1$ . Let: T be a set of ordinals  $\langle \alpha \rangle$  such that  $\alpha = \operatorname{ssup} T$
    - $\langle 3 \rangle 2$ .  $\alpha = \sup T$ 
      - $\langle 4 \rangle 1$ . For all  $\beta \in T$  we have  $\beta \leq \alpha$
      - $\langle 4 \rangle 2$ . Let:  $\mu$  be any upper bound for T Prove:  $\alpha \leq \mu$
      - $\langle 4 \rangle 3. \ \alpha \leq \mu + 1$

PROOF: Since  $\forall \beta \in T.\beta < \mu + 1$ .

 $\langle 4 \rangle 4$ .  $\alpha \neq \mu + 1$ 

PROOF: Since  $\alpha$  is a limit ordinal.

- $\langle 4 \rangle 5$ .  $\alpha < \mu + 1$
- $\langle 4 \rangle 6. \ \alpha \leq \mu$
- $\langle 3 \rangle 3$ . cf  $\alpha \leq |T|$

### 12.5 Inaccessible Cardinals

**Definition 12.5.1** (Inaccessible Cardinal). A cardinal number  $\kappa$  is *inaccessible* iff:

- $\kappa > \aleph_0$
- $\forall \lambda < \kappa.2^{\lambda} < \kappa$  (cardinal exponentiation)
- $\kappa$  is regular.

Any inaccessible cardinal is weakly inaccessible.

### Proof:

- $\langle 1 \rangle 1$ . Let:  $\kappa = \aleph_{\lambda}$  be weakly inaccessible. Prove:  $\lambda$  is a limit ordinal.
- $\langle 1 \rangle 2. \ \lambda \neq 0$
- $\langle 1 \rangle 3$ . Assume: for a contradiction  $\lambda = \beta + 1$
- $\langle 1 \rangle 4$ .  $\aleph_{\beta} < \kappa$
- $\langle 1 \rangle 5. \ 2^{\aleph_{\beta}} < \kappa$
- $\langle 1 \rangle 6. \ \aleph_{\beta+1} < \kappa$

PROOF: Since  $\aleph_{\beta+1} \leq 2^{\aleph_{\beta}}$ .

 $\langle 1 \rangle 7$ . Q.E.D.

PROOF: This contradicts  $\langle 1 \rangle 3$ .

**Proposition 12.5.2.** If the Generalized Continuum Hypothesis is true, then every weakly inaccessible cardinal is inaccessible.

### Proof:

(1)1. Assume: The Generalized Continuum Hypothesis.

 $= \kappa$ 

- $\langle 1 \rangle 2$ . Let:  $\kappa = \aleph_{\lambda}$  be weakly inaccessible.
- $\langle 1 \rangle 3. \ \kappa > \aleph_0$

PROOF:  $\lambda > 0$  because  $\lambda$  is a limit ordinal.

- $\langle 1 \rangle 4$ . For all  $\mu < \kappa$  we have  $2^{\mu} < \kappa$ 
  - $\langle 2 \rangle 1$ . Let:  $\mu < \kappa$
  - $\langle 2 \rangle 2$ . Let:  $\mu = \aleph_{\alpha}$
  - $\langle 2 \rangle 3. \ \alpha < \lambda$
  - $\langle 2 \rangle 4$ .  $\alpha + 1 < \lambda$

PROOF:  $\lambda$  is a limit ordinal.

 $\langle 2 \rangle 5$ .  $2^{\mu} < \kappa$ 

Proof:

$$\begin{aligned} 2^{\mu} &= 2^{\aleph_{\alpha}} & (\langle 2 \rangle 2) \\ &= 2^{\beth_{\alpha}} & (\langle 1 \rangle 1) \\ &= \beth_{\alpha+1} & (\langle 1 \rangle 1) \\ &= \aleph_{\alpha+1} & (\langle 1 \rangle 1) \\ &< \aleph_{\lambda} & (\langle 2 \rangle 4) \end{aligned}$$

 $(\langle 1 \rangle 2)$ 

 $\langle 1 \rangle$ 5.  $\kappa$  is regular. PROOF:  $\langle 1 \rangle$ 2

**Lemma 12.5.3.** Let  $\kappa$  be an inaccessible cardinal. For every ordinal  $\alpha < \kappa$  we have  $\beth_{\alpha} < \kappa$ .

Proof:

 $\langle 1 \rangle 1. \ \ \beth_0 < \kappa$ 

PROOF: Since  $\kappa > \aleph_0$ .

 $\langle 1 \rangle 2$ . For any ordinal  $\alpha$ , if  $\beth_{\alpha} < \kappa$  then  $\beth_{\alpha+1} < \kappa$ .

PROOF: Since  $\beth_{\alpha+1} = 2^{\beth_{\alpha}} < \kappa$ .

 $\langle 1 \rangle 3$ . For any limit ordinal  $\lambda$ , if  $\forall \alpha < \lambda. \beth_{\alpha} < \kappa$  and  $\lambda < \kappa$  then  $\beth_{\lambda} < \kappa$ .

PROOF: By regularity of  $\kappa$ , since  $\beth_{\lambda}$  is the union of  $|\lambda|$  cardinals all  $< \kappa$ .

**Lemma 12.5.4.** Let  $\kappa$  be an inaccessible cardinal. For all  $A \in V_{\kappa}$  we have  $|A| < \kappa$ .

Proof:

 $\langle 1 \rangle 1$ . Let:  $A \in V_{\kappa}$ 

 $\langle 1 \rangle 2$ . PICK  $\alpha < \kappa$  such that  $A \in V_{\alpha}$ 

 $\langle 1 \rangle 3. \ A \subseteq V_{\alpha}$ 

 $\langle 1 \rangle 4. \ |A| \leq |V_{\alpha}| \leq \beth_{\alpha} < \kappa$ 

**Theorem Schema 12.5.5.** For every axiom  $\alpha$  of ZFC, the following is a theorem:

For any inaccessible cardinal  $\kappa$ , we have  $V_{\kappa}$  is a model of  $\alpha$ .

PROOF: For every axiom except the Replacement Axioms, we have Corollary 11.0.11.1.

For an Axiom of Replacement using the property  $P[x, y, z_1, \dots, z_n]$ , we reason as follows:

 $\langle 1 \rangle 1$ . Let:  $\kappa$  be an inaccessible cardinal

PROVE:

$$\forall a_1, \dots, a_n, B \in V_{\kappa} (\forall x \in B. \forall y, y' \in V_{\kappa} \\ (P[x, y, a_1, \dots, a_n]^{V_{\kappa}} \land P[x, y', a_1, \dots, a_n]^{V_{\kappa}} \Rightarrow y = y') \Rightarrow \\ \exists C \in V_{\kappa} \forall y \in V_{\kappa} (y \in C \Leftrightarrow \exists x \in B. P[x, y, a_1, \dots, a_n]^{V_{\kappa}}))$$

 $\langle 1 \rangle 2$ . Let:  $a_1, \ldots, a_n, B \in V_{\kappa}$ 

 $\langle 1 \rangle 3$ . Assume: for all  $x \in B$ , there exists at most one  $y \in V_{\kappa}$  such that  $P[x,y,a_1,\ldots,a_n]^{V_{\kappa}}$ .

 $\langle 1 \rangle 4$ . Let:  $F = \{(x, y) \in B \times V_{\kappa} \mid P[x, y, a_1, \dots, a_n]^{V_{\kappa}} \}$ 

 $\langle 1 \rangle 5$ . Let:  $C = \operatorname{ran} F$ 

Prove:  $C \in V_{\kappa}$ 

 $\langle 1 \rangle 6$ . Let:  $S = \{ \operatorname{rank} F(x) \mid x \in \operatorname{dom} F \}$ 

 $\langle 1 \rangle 7$ .  $|S| < \kappa$ 

PROOF: Since  $|S| \leq |\operatorname{dom} F| \leq |B| < \kappa$ .

```
\begin{split} &\langle 1 \rangle 8. \  \, \forall \alpha \in S.\alpha < \kappa \\ & \text{Proof: Since } F(x) \in V_\kappa \text{ for all } x \in \operatorname{dom} F. \\ &\langle 1 \rangle 9. \ \sup S < \kappa \\ & \text{Proof: Since } \kappa \text{ is regular.} \\ &\langle 1 \rangle 10. \ \operatorname{rank} C \leq \sup S + 1 \\ &\langle 1 \rangle 11. \ \operatorname{rank} C < \kappa \\ &\langle 1 \rangle 12. \  \, C \in V_\kappa \\ & \Box \end{split}
```

# Chapter 13

# Group Theory

## 13.1 Groups

**Definition 13.1.1** (Group). A group G consists of a set G and a function  $\cdot: G^2 \to G$  such that:

- $1. \cdot is associative$
- 2. There exists  $e \in G$  such that  $\forall x \in G.xe = x$  and  $\forall x \in G.\exists y \in G.xy = e$ .

Proposition 13.1.2. The inverse of an element in a group is unique.

Proof:

 $\langle 1 \rangle 1$ . Assume: b and b' are inverses of a.

 $\langle 1 \rangle 2. \ b = b'$ 

Proof:

$$b = be$$

$$= bab'$$

$$= eb'$$

$$= b'$$

**Definition 13.1.3.** We write  $x^{-1}$  for the inverse of x.

**Proposition 13.1.4.** In any group, if ab = ac then b = c.

Proof:

$$b = eb$$

$$= a^{-1}ab$$

$$= a^{-1}ac$$

$$= ec$$

$$= c$$

**Definition 13.1.5.** Let Grp be the category of groups and group homomorphisms.

**Definition 13.1.6.** We identify any group G with the category G with one object whose morphisms are the elements of G, with composition given by the multiplication in G.

## 13.2 Abelian Groups

**Definition 13.2.1** (Abelian group). An *Abelian group* is a group whose multiplication is commutative.

We may say we are writing an Abelian group *additively*, meaning we write a + b for ab, 0 for e and -a for  $a^{-1}$ . In this case we write a - b for  $ab^{-1}$ .

## Chapter 14

# Ring Theory

## 14.1 Rings

**Definition 14.1.1** (Commutative Ring). A *commutative ring* consists of a set R and two binary operations +,  $\cdot$  on R such that:

- D is an Abelian group under +. Let us write 0 for its identity element.
- $\bullet$  · is commutative and associative, and distributes over +.
- $\bullet$  · has an identity element 1 that is different from 0.

**Proposition 14.1.2.** In any commutative ring, 0x = 0.

Proof:

$$(0+0)x = 0x$$
∴  $0x + 0x = 0x + 0$ 
∴  $0x = 0$  (Proposition 13.1.4)

**Proposition 14.1.3.** In any commutative ring, (-a)b = -(ab).

Proof:

$$ab + (-a)b = (a + (-a))b$$
  
=  $0b$   
=  $0$  (Proposition 14.1.2) $\square$ 

## 14.2 Ordered Rings

**Definition 14.2.1** (Ordered Commutative Ring). An ordered commutative ring consists of a commutative ring R with a linear order < on R such that:

• for all  $x, y, z \in R$ , we have x < y if and only if x + z < y + z.

• for all  $x, y, z \in R$ , if 0 < z then we have x < y if and only if xz < yz.

**Proposition 14.2.2.** In any ordered commutative ring, 0 < 1.

PROOF: If 1 < 0 then we have 0 < -1 and so 0 < (-1)(-1) = 1, which is a contradiction.  $\square$ 

**Proposition 14.2.3.** The ordering on an ordered commutative ring is dense; that is, if x < y then there exists z such that x < z < y.

PROOF: Take z = (x + y)/2.  $\square$ 

## 14.3 Integral Domains

**Definition 14.3.1** (Integral Domain). An *integral domain* is a commutative ring such that, for all  $a, b \in D$ , if ab = 0 then a = 0 or b = 0.

**Proposition 14.3.2.** In any integral domain, if ab = ac and  $a \neq 0$  then b = c.

PROOF: We have a(b-c)=0 and  $a\neq 0$  so b-c=0 hence b=c.  $\square$ 

**Definition 14.3.3** (Ordered Integral Domain). An *ordered integral domain* is an ordered commutative ring that is an integral domain.

## Chapter 15

# Field Theory

### 15.1 Fields

**Definition 15.1.1** (Field). A *field* F is a commutative ring such that  $0 \neq 1$  and, for all  $x \in F$ , if  $x \neq 0$  then there exists  $y \in F$  such that xy = 1.

Proposition 15.1.2. Every field is an integral domain.

PROOF: If ab = 0 and  $a \neq 0$  then  $b = a^{-1}ab = 0$ .

**Proposition 15.1.3.** In any field F, we have  $F - \{0\}$  is an Abelian group under multiplication.

PROOF: Immediate from the definition.  $\Box$ 

**Definition 15.1.4** (Field of Fractions). Let D be an integral domain. The *field of fractions* of D is the quotient set  $F = (D \times (D - \{0\})) / \sim$  where

$$(a,b) \sim (c,d) \Leftrightarrow ad = bc$$

under

$$[(a,b)] + [(c,d)] = [(ad+bc,bd)]$$
$$[(a,b)][(c,d)] = [(ac,bd)]$$

We prove this is a field.

Proof:

 $\langle 1 \rangle 1$ .  $\sim$  is an equivalence relation on  $D \times (D - \{0\})$ . PROOF:

 $\langle 2 \rangle 1$ .  $\sim$  is reflexive.

PROOF: We always have ab = ba.

 $\langle 2 \rangle 2$ . ~ is symmetric.

PROOF: If ad = bc then cb = da.

```
\langle 2 \rangle 3. \sim is transitive.
     \langle 3 \rangle 1. Assume: (a,b) \sim (c,d) \sim (e,f)
     \langle 3 \rangle 2. ad = bc and cf = de
     \langle 3 \rangle 3. adf = bde
        PROOF: adf = bcf = bde
     \langle 3 \rangle 4. af = be
        Proof: Proposition 14.3.2.
\langle 1 \rangle 2. Addition is well-defined.
  Proof:
  \langle 2 \rangle 1. If b \neq 0 and d \neq 0 then bd \neq 0.
     PROOF: Since D is an integral domain.
  \langle 2 \rangle 2. If ab' = a'b and cd' = c'd then (ad + bc)b'd' = (a'd' + b'c')bd.
     Proof:
                                 (ad + bc)b'd' = ab'dd' + bb'cd'
                                                 = a'bdd' + bb'c'd
                                                 = (a'd' + b'c')bd
\langle 1 \rangle 3. Multiplication is well-defined.
  Proof:
  \langle 2 \rangle 1. If b \neq 0 and d \neq 0 then bd \neq 0.
     PROOF: Since D is an integral domain.
  (2)2. If [(a,b)] = [(a',b')] and [(c,d)] = [(c',d')] then [(ac,bd)] = [(a'c',b'd')].
     PROOF: If ab' = a'b and cd' = c'd then acb'd' = a'c'bd.
\langle 1 \rangle 4. Addition is commutative.
  PROOF: [(a,b)] + [(c,d)] = [(ad+bc,bd)] = [(cb+da,db)] = [(c,d)] + [(a,b)] \sqcup
\langle 1 \rangle5. Addition is associative.
  Proof:
          [(a,b)] + ([(c,d)] + [(e,f)]) = [(a,b)] + [(cf + de, df)]
                                            = [(adf + bcf + bde, bdf)]
                                            = [(ad + bc, bd)] + [(e, f)]
                                            = ([(a,b)] + [(c,d)]) + [(e,f)]
                                                                                       \langle 1 \rangle 6. For any x \in F we have x + [(0,1)] = x
  PROOF: [(a,b)] + [(0,1)] = [(a \cdot 1 + b \cdot 0, b \cdot 1)] = [(a,b)] \sqcup
\langle 1 \rangle 7. For any x \in F, there exists y \in F such that x + y = [(0,1)].
  PROOF: [(a,b)] + [(-a,b)] = [(ab-ab,b^2)] = [(0,b^2)] = [(0,1)]
\langle 1 \rangle 8. Multiplication is commutative.
  PROOF: [(a,b)][(c,d)] = [(c,d)][(a,b)] = [(ac,bd)].
\langle 1 \rangle 9. Multiplication is assocative.
  PROOF: [(a,b)]([(c,d)][(e,f)]) = ([(a,b)][(c,d)])[(e,f)] = [(ace,bdf)].
\langle 1 \rangle 10. For any x \in F we have x[(1,1)] = x
  PROOF: [(a,b)][(1,1)] = [(a,b)]
\langle 1 \rangle 11. For any non-zero x \in F, there exists y \in F such that xy = [(1,1)].
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Proof:
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 \begin{array}{l} \langle 2 \rangle 1. \ \ \mathrm{Let:} \ \left[ (a,b) \right] \in \mathbb{Q} \\ \langle 2 \rangle 2. \ \ \mathrm{Assume:} \ \left[ (a,b) \right] \neq \left[ (0,1) \right] \\ \langle 2 \rangle 3. \ \ a \neq 0 \\ \langle 2 \rangle 4. \ \ \left[ (a,b) \right] \left[ (b,a) \right] = \left[ (1,1) \right] \\ \square \\ \end{array}
```

**Definition 15.1.5.** For any field F, let N(F) be the intersection of all the subsets  $S \subseteq F$  such that  $1 \in S$  and  $\forall x \in S.x + 1 \in S$ .

**Definition 15.1.6** (Characteristic Zero). A field F has *characteristic* 0 iff  $0 \notin N(F)$ .

**Proposition 15.1.7.** In a field F with characteristic 0, the function  $n: \mathbb{N} \to N(F)$  defined by

$$n(0) = 1$$
$$n(x+1) = n(x) + 1$$

is a bijection.

### Proof:

 $\langle 1 \rangle 1$ . *n* is injective.

 $\langle 2 \rangle 1$ . Assume: for a contradiction n(i) = n(j) with  $i \neq j$ 

 $\langle 2 \rangle 2$ . Assume: w.l.o.g. i < j

 $\langle 2 \rangle 3$ . n(j-i)=0

 $\langle 2 \rangle 4$ . Q.E.D.

PROOF: This contradicts the fact that F has characteristic 0.

 $\langle 1 \rangle 2$ . n is surjective.

PROOF: Since ran n is a subset of F that includes 1 and is closed under +1.

**Definition 15.1.8.** In any field F, let

$$I(F) = N(F) \cup \{0\} \cup \{-x \mid x \in N(F)\}\$$

**Definition 15.1.9.** In any field F, let

$$Q(F) = \{x/y \mid x, y \in I(F), y \neq 0\}$$

**Proposition 15.1.10.** Q(F) is the smallest subfield of F.

PROOF: Q(F) is closed under + and  $\cdot$ , and any subset of F closed under + and  $\cdot$  that contains 0 and 1 must include Q(F).  $\square$ 

**Theorem 15.1.11.** Let F and G be fields of characteristic 0. Then there exists a unique field isomorphism between Q(F) and Q(G).

- $\langle 1 \rangle 1$ . Let:  $\phi: N(F) \to N(G)$  be the unique function such that  $\phi(1) = 1$  and  $\forall x \in N(F). \phi(x+1) = \phi(x) + 1$ .
- $\langle 1 \rangle 2$ .  $\phi$  is a bijection.

Proof: Similar to Proposition 15.1.7.

 $\langle 1 \rangle 3. \ \forall x, y \in N(F). \phi(x+y) = \phi(x) + \phi(y)$ 

PROOF: Induction on y.

 $\langle 1 \rangle 4. \ \forall x, y \in N(F). \phi(xy) = \phi(x)\phi(y)$ 

PROOF: Induction on y.

- (1)5. Extend  $\phi$  to a bijection  $I(F) \cong I(G)$  such that  $\forall x, y \in I(F).\phi(x+y) = \phi(x) + \phi(y)$  and  $\forall x, y \in I(F).\phi(xy) = \phi(x)\phi(y)$ 
  - $\langle 2 \rangle 1$ . Define  $\phi(0) = 0$  and  $\phi(-x) = -\phi(x)$  for  $x \in N(F)$ 
    - $\langle 3 \rangle 1. \ 0 \notin N(F)$
    - $\langle 3 \rangle 2$ . For all  $x \in N(F)$  we have  $-x \notin N(F)$

PROOF: Then we would have  $x + -x = 0 \in N(F)$ .

- $\langle 3 \rangle 3$ . For all  $x \in N(F)$  we have  $-x \neq 0$
- $\langle 2 \rangle 2$ . For all  $x, y \in I(F)$  we have  $\phi(x+y) = \phi(x) + \phi(y)$

PROOF: Case analysis on x and y.

 $\langle 2 \rangle 3$ . For all  $x, y \in I(F)$  we have  $\phi(xy) = \phi(x)\phi(y)$ 

PROOF: Case analysis on x and y.

- $\langle 1 \rangle$ 6. Extend  $\phi$  to a bijection  $Q(F) \cong Q(G)$  such that  $\forall x, y \in Q(F).\phi(x+y) = \phi(x) + \phi(y)$  and  $\forall x, y \in Q(F).\phi(xy) = \phi(x)\phi(y)$ 
  - $\langle 2 \rangle 1$ . Define  $\phi(x/y) = \phi(x)/\phi(y)$
- $\langle 1 \rangle 7$ .  $\phi$  is unique.
  - $\langle 2 \rangle 1$ . Let:  $\theta$  satisfy the theorem.
  - $\langle 2 \rangle 2$ . For all  $x \in N(F)$  we have  $\theta(x) = \phi(x)$
  - $\langle 2 \rangle 3$ . For all  $x \in I(F)$  we have  $\theta(x) = \phi(x)$
  - $\langle 2 \rangle 4$ . For all  $x \in Q(F)$  we have  $\theta(x) = \phi(x)$

## 15.2 Ordered Fields

**Definition 15.2.1** (Ordered Field). An *ordered field* is an ordered commutative ring that is a field.

**Proposition 15.2.2.** Every ordered field F has characteristic  $\theta$ .

PROOF: We have 0 < n for all  $n \in N(F)$ .  $\square$ 

**Proposition 15.2.3.** Let F be a field of characteristic 0. Then there exists a unique relation < on Q(F) that makes Q(F) into an ordered field.

Proof: Easy.  $\square$ 

Corollary 15.2.3.1. Let F and G be ordered fields. Let  $\phi$  be the unique field isomorphism between Q(F) and Q(G). Then  $\phi$  is an ordered field isomorphism.

**Definition 15.2.4** (Archimedean). An ordered field F is Archimedean iff

$$\forall x \in F. \exists n \in N(F). n > x .$$

**Proposition 15.2.5.** Let F be an Archimedean ordered field. Let  $x, y \in F$  with x > 0. Then there exists  $n \in N(F)$  such that nx > y.

PROOF: Pick n > y/x.  $\square$ 

**Proposition 15.2.6.** Let F be an Archimedean ordered field. For all  $x, y \in F$ , if x < y, then there exists  $r \in Q(F)$  such that x < r < y.

### Proof:

- $\langle 1 \rangle 1$ . Case: x > 0
  - $\langle 2 \rangle 1$ . PICK  $n \in N(F)$  such that n(y-x) > 1

Proof: Proposition 15.2.5.

- $\langle 2 \rangle 2$ . ny > 1 + nx
- $\langle 2 \rangle$ 3. Let: m be the least element of N(F) such that m > nx.
- $\langle 2 \rangle 4$ .  $m-1 \leq nx$
- $\langle 2 \rangle 5$ . nx < m < ny
- $\langle 2 \rangle 6$ . x < m/n < y
- $\langle 1 \rangle 2$ . Case:  $x \leq 0$ 
  - $\langle 2 \rangle 1$ . PICK  $k \in N(F)$  such that k > -x
  - $\langle 2 \rangle 2$ . 0 < x + k < y + k
  - $\langle 2 \rangle 3$ . Pick  $r \in Q(F)$  such that x + k < r < y + k

Proof:  $\langle 1 \rangle 1$ 

 $\langle 2 \rangle 4$ . x < r - k < y

**Definition 15.2.7** (Complete). An ordered field F is *complete* iff every nonempty subset of F bounded above has a least upper bound.

**Proposition 15.2.8.** Every complete ordered field is Archimedean.

#### Proof:

- $\langle 1 \rangle 1$ . Let: F be a complete ordered field.
- $\langle 1 \rangle 2$ . Let:  $x \in F$
- $\langle 1 \rangle 3$ . Assume: for a contradiction there is no member of N(F) greater than x.
- $\langle 1 \rangle 4$ . x is an upper bound for N(F).
- $\langle 1 \rangle 5$ . Let:  $y = \sup N(F)$
- $\langle 1 \rangle 6$ . Pick  $n \in N(F)$  such that y 1 < n
- $\langle 1 \rangle 7$ . y < n+1
- $\langle 1 \rangle 8$ . Q.E.D.

Proof: This is a contradiction.

**Proposition 15.2.9.** Let F be a complete ordered field and  $a \in F$  be nonnegative. Then there exists  $b \in F$  such that  $b^2 = a$ .

- $\langle 1 \rangle 1$ . Let:  $B = \{ x \in F \mid 0 \le x \le 1 + a \}$
- $\langle 1 \rangle 2$ . Let:  $\phi : B \to B$  be the function

$$\phi(x) = x + \frac{1}{2(1+a)}(a-x^2) .$$

- $\langle 1 \rangle 3$ .  $\phi$  is strictly monotone.
  - $\langle 2 \rangle$ 1. Let:  $0 \le x < y \le 1 + a$  $\langle 2 \rangle$ 2.  $1 \frac{x+y}{2(1+a)} > 0$

  - $\langle 2 \rangle 3. \ \phi(y) \phi(x) = (y x)(1 \frac{x+y}{2(1+a)}) > 0$
  - $\langle 2 \rangle 4. \ \phi(x) < \phi(y)$
- $\langle 1 \rangle 4$ . Pick  $b \in B$  such that  $\phi(b) = b$ .

PROOF: Knaster Fixed-Point Theorem.

$$\langle 1 \rangle 5. \ b^2 = a$$

**Theorem 15.2.10** (Uniqueness of the Complete Ordered Field). If F and G are complete ordered fields, then there exists a unique bijection  $\phi: F \cong G$  such that, for all  $x, y \in F$ ,

$$\phi(x+y) = \phi(x) + \phi(y)$$
$$\phi(xy) = \phi(x)\phi(y)$$

This bijection also satisfies: for all  $x, y \in F$ ,

$$x < y \Leftrightarrow \phi(x) < \phi(y)$$
.

#### Proof:

 $\langle 1 \rangle 1$ . Pick a bijection  $\phi: Q(F) \cong Q(G)$  such that, for all  $x, y \in Q(F)$ ,

$$\phi(x+y) = \phi(x) + \phi(y)$$
$$\phi(xy) = \phi(x)\phi(y)$$

$$x < y \Leftrightarrow \phi(x) < \phi(y)$$

Proof: Corollary 15.2.3.1.

 $\langle 1 \rangle 2$ . Q(F) intersects every interval in F.

Proof: Proposition 15.2.6.

 $\langle 1 \rangle 3$ . Q(G) intersects every interval in G.

Proof: Proposition 15.2.6.

 $\langle 1 \rangle 4$ . PICK an order isomorphism  $\psi : F \cong G$  that extends  $\phi$ .

PROOF: Theorem 6.1.21.

- $\langle 1 \rangle 5. \ \forall x, y \in F. \psi(x+y) = \psi(x) + \psi(y)$ 
  - $\langle 2 \rangle 1$ . Let:  $x, y \in F$
  - $\langle 2 \rangle 2$ .  $\psi(x) + \psi(y) \not< \psi(x+y)$ 
    - $\langle 3 \rangle 1$ . Assume: for a contradiction  $\psi(x) + \psi(y) < \psi(x+y)$
    - $\langle 3 \rangle 2$ . Pick  $r' \in Q(G)$  such that  $\psi(x) < r' < \psi(x+y) \psi(y)$
    - $\langle 3 \rangle 3$ . Pick  $s' \in Q(G)$  such that  $\psi(y) < s' < \psi(x+y) r'$
    - $\langle 3 \rangle 4. \ r' + s' < \psi(x+y)$
    - $\langle 3 \rangle 5$ . Pick  $r, s \in Q(F)$  such that  $\phi(r) = r'$  and  $\phi(s) = s'$
    - $\langle 3 \rangle 6. \ \phi(r+s) = r' + s'$
    - $\langle 3 \rangle 7. \ \psi(x) < \psi(r)$
    - $\langle 3 \rangle 8. \ \psi(y) < \psi(s)$
    - $\langle 3 \rangle 9. \ \psi(x+y) > \psi(r+s)$
    - $\langle 3 \rangle 10. \ x < r$

```
\langle 3 \rangle 11. \ y < s
       \langle 3 \rangle 12. x + y > r + s
       \langle 3 \rangle 13. Q.E.D.
           PROOF: This is a contradiction.
   \langle 2 \rangle 3. \ \psi(x+y) \not< \psi(x) + \psi(y)
       Proof: Similar.
\langle 1 \rangle 6. \ \forall x, y \in F. \psi(xy) = \psi(x) \psi(y)
    \langle 2 \rangle 1. Let: x, y \in F
   \langle 2 \rangle 2. Case: x and y are positive.
       \langle 3 \rangle 1. \ \psi(x)\psi(y) \not< \psi(xy)
           \langle 4 \rangle1. Assume: for a contradiction \psi(x)\psi(y) < \psi(xy)
           \langle 4 \rangle 2. PICK r' \in Q(G) such that \psi(x) < r' < \psi(xy)/\psi(y)
           \langle 4 \rangle 3. Pick s' \in Q(G) such that \psi(y) < s' < \psi(xy)/r'
           \langle 4 \rangle 4. r's' < \psi(xy)
           \langle 4 \rangle5. PICK r, s \in Q(F) such that \phi(r) = r' and \phi(s) = s'
           \langle 4 \rangle 6. \ \phi(rs) = r's'
           \langle 4 \rangle 7. x < r, y < s \text{ and } rs < xy
           \langle 4 \rangle 8. Q.E.D.
              PROOF: This is a contradiction.
       \langle 3 \rangle 2. \ \psi(xy) \not< \psi(x)\psi(y)
           PROOF: Similar.
   \langle 2 \rangle 3. Case: x and y are not both positive.
       PROOF: Follows from \langle 2 \rangle 2 since \psi(-x) = -\psi(x) by \langle 1 \rangle 5.
\langle 1 \rangle 7. For any field isomorphism \theta : F \cong G, we have \theta = \psi.
   \langle 2 \rangle 1. \ \theta \upharpoonright Q(F) = \phi
       PROOF: Theorem 15.1.11.
   \langle 2 \rangle 2. \theta is strictly monotone.
       \langle 3 \rangle 1. Let: x, y \in F with x < y
       \langle 3 \rangle 2. y - x > 0
       \langle 3 \rangle 3. PICK z \in F such that z^2 = y - x
       \langle 3 \rangle 4. \theta(z)^2 = \theta(y) - \theta(x)
       \langle 3 \rangle 5. \theta(y) - \theta(x) > 0
       \langle 3 \rangle 6. \ \theta(x) < \theta(y)
   \langle 2 \rangle 3. \ \theta = \psi
       Proof: By the uniqueness of \psi.
```

## Chapter 16

# Number Systems

## 16.1 The Integers

**Definition 16.1.1.** The set of integers  $\mathbb{Z}$  is the quotient set  $\mathbb{N}^2/\sim$ , where  $(m,n)\sim(p,q)$  iff m+q=n+p.

We prove  $\sim$  is an equivalence relation on  $\mathbb{N}^2$ .

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Proof:
```

 $\langle 1 \rangle 1$ .  $\sim$  is reflexive.

PROOF: For all  $m, n \in \mathbb{N}$  we have m + n = n + m.

 $\langle 1 \rangle 2$ .  $\sim$  is symmetric.

PROOF: If m + q = n + p then p + n = q + m.

 $\langle 1 \rangle 3$ .  $\sim$  is transitive.

- $\langle 2 \rangle 1$ . Assume:  $(m,n) \sim (p,q) \sim (r,s)$
- $\langle 2 \rangle 2$ . m+q=n+p and p+s=q+r
- $\langle 2 \rangle 3$ . m+q+s=n+q+r
- $\langle 2 \rangle 4$ . m+s=n+r

Proof: By cancellation.

**Definition 16.1.2** (Addition). Define  $addition + \text{ on } \mathbb{Z}$  by [(m,n)] + [(p,q)] = [(m+p,n+q)].

We prove this is well-defined.

PROOF: If m+n'=n+m' and p+q'=q+p' then m+p+n'+q'=n+q+m'+p'.

**Proposition 16.1.3.** Addition on  $\mathbb{Z}$  is commutative.

PROOF: 
$$[(m,n)] + [(p,q)] = [(m+p,n+q)] = [(p+m,q+n)] = [(p,q)] + [(m,n)].$$

**Proposition 16.1.4.** Addition on  $\mathbb{Z}$  is associative.

PROOF: [(m,n)] + ([(p,q)] + [(r,s)]) = [(m+p+r,n+q+s)] = ([(m,n)] + [(p,q)]) + [(r,s)].

**Proposition 16.1.5.** Given natural numbers m and n, we have [(m,0)] = [(n,0)] iff m = n.

PROOF: Immediate from definitions.

**Definition 16.1.6.** We identify any natural number n with the integer [(n,0)].

**Proposition 16.1.7.** Addition on integers agrees with addition on natural numbers.

PROOF: Since [(m,0)] + [(n,0)] = [(m+n,0)].

**Proposition 16.1.8.** For all  $a \in \mathbb{Z}$  we have a + 0 = a.

PROOF: [(m, n)] + [(0, 0)] = [(m + 0, n + 0)] = [(m, n)].

**Proposition 16.1.9.** For all  $a \in \mathbb{Z}$ , there exists  $b \in \mathbb{Z}$  such that a + b = 0.

PROOF: [(m,n)] + [(n,m)] = [(m+n,m+n)] = [(0,0)]

Proposition 16.1.10. The integers form an Abelian group under addition.

Proof: Proposition 16.1.3, 16.1.4, 16.1.8, 16.1.9.

**Definition 16.1.11.** Define multiplication  $\cdot$  on  $\mathbb{Z}$  by: [(m,n)][(p,q)] = [(mp + nq, mq + np)].

We prove this is well defined.

### Proof:

- $\langle 1 \rangle 1.$  Assume: m+n'=n+m' and p+q'=q+p' Prove: mp+nq+m'q'+n'p'=mq+np+m'p'+n'q'
- $\langle 1 \rangle 2$ . mp + n'p = np + m'p
- $\langle 1 \rangle 3$ . nq + m'q = mq + n'q
- $\langle 1 \rangle 4$ . m'p + m'q' = m'q + m'p'
- $\langle 1 \rangle 5. \ n'q + n'p' = n'p + n'q'$
- (1)6. mp + n'p + nq + m'q + m'p + m'q' + n'q + n'p' = np + m'p + mq + n'q + m'q + m'p' + n'p + n'q'
- $\langle 1 \rangle 7$ . mp + nq + m'q' + n'p' = mq + np + m'p' + n'q'

Proof: By cancellation.

**Proposition 16.1.12.** Multiplication on integers agrees with multiplication on natural numbers.

PROOF: Since [(m,0)][(n,0)] = [(mn+0,m0+n0)] = [(mn,0)].

**Proposition 16.1.13.** *Multiplication on*  $\mathbb{Z}$  *is commutative.* 

PROOF: [(m,n)][(p,q)] = [(mp+nq,mq+np)] = [(pm+qn,pn+qm)] = [(p,q)][(m,n)].

**Proposition 16.1.14.** *Multiplication on*  $\mathbb{Z}$  *is associative.* 

Proof:

$$\begin{split} [(m,n)]([(p,q)][(r,s)]) &= [(m,n)][(pr+qs,ps+qr)] \\ &= [(mpr+mqs+nps+nqr,mps+mqr+npr+nqs)] \\ &= [(mp+nq,mq+np)][(r,s)] \\ &= ([(m,n)][(p,q)])[(r,s)] \end{split}$$

Proposition 16.1.15. Multiplication distributes over addition.

Proof:

$$\begin{split} [(m,n)]([(p,q)]+[(r,s)]) &= [(m,n)][(p+r,q+s)] \\ &= [(mp+mr+nq+ns,np+nr+mq+ms)] \\ [(m,n)][(p,q)]+[(m,n)][(r,s)] &= [(mp+nq,mq+np)]+[(mr+ns,ms+nr)] \\ &= [(mp+nq+mr+ns,mq+np+ms+nr)] \end{split}$$

**Proposition 16.1.16.** For any integer a we have a1 = a.

PROOF: Since 
$$[(m,n)][(1,0)] = [(m1+n0,m0+n1)] = [(m,n)]$$
.

**Proposition 16.1.17.** For any integers a and b, if ab = 0 then a = 0 or b = 0.

Proof:

```
\langle 1 \rangle 1. Assume: [(m,n)][(p,q)] = [(0,0)]
\langle 1 \rangle 2. mp + nq = mq + np
\langle 1 \rangle 3. Assume: [(m,n)] \neq [(0,0)]
\langle 1 \rangle 4. m \neq n
```

PROVE: p = q

 $\langle 1 \rangle 5$ . Case: m < n

 $\langle 2 \rangle 1. \ p \not< q$ 

PROOF: If p < q then mq + np < mp + nq by Proposition 9.4.6.

 $\langle 2 \rangle 2$ .  $q \not< p$ 

PROOF: If q < p then mp + nq < mq + np by Proposition 9.4.6.

 $\langle 2 \rangle 3. \ p = q$ 

PROOF: By trichotomy.

 $\langle 1 \rangle 6$ . Case: n < m

PROOF: Similar.

**Proposition 16.1.18.** The integers  $\mathbb{Z}$  form an integral domain.

PROOF: Propositions 16.1.13, 16.1.14, 16.1.15, 16.1.16, 16.1.17, 16.1.10.

**Definition 16.1.19.** Define < on  $\mathbb{Z}$  by [(m,n)] < [(p,q)] if and only if m+q < n+p.

We prove this is well-defined.

$$\begin{array}{ll} \langle 1 \rangle 1. & \text{Assume: } m+n'=n+m' \text{ and } p+q'=q+p'. \\ & \text{Prove: } m+q< n+p \text{ if and only if } m'+q'< n'+p' \\ \langle 1 \rangle 2. & m+q< n+p \text{ if and only if } m'+q'< n'+p' \\ & \text{Proof: } \\ & m+q< n+p \Leftrightarrow m+n'+q< n+n'+p \\ & \Leftrightarrow m'+n+q< n+n'+p \\ & \Leftrightarrow m'+q< n'+p \\ & \Leftrightarrow m'+q+p'< n'+p+p' \end{array} \qquad \begin{array}{l} \text{(Corollary 7.7.7.1)} \\ & \Leftrightarrow m'+q+p'< n'+p+p' \\ & \Leftrightarrow m'+q'+p+p' \end{array}$$

**Proposition 16.1.20.** The ordering on the integers agrees with the ordering on the natural numbers.

PROOF: We have [(m,0)] < [(n,0)] iff m < n.  $\square$ 

**Proposition 16.1.21.** < is a linear order on  $\mathbb{Z}$ .

### Proof:

 $\langle 1 \rangle 1$ . < is irreflexive.

PROOF: We never have m + n < m + n.

- $\langle 1 \rangle 2$ . < is transitive.
  - $\langle 2 \rangle 1$ . Assume: [(m,n)] < [(p,q)] < [(r,s)]
  - $\langle 2 \rangle 2$ . m+q < n+p and p+s < q+r
  - $\langle 2 \rangle 3. \ m + q + s < n + q + r$

PROOF: m + q + s < n + p + s < n + q + r

 $\langle 2 \rangle 4$ . m+s < n+r

PROOF: Corollary 7.7.7.1.

 $\langle 1 \rangle 3.$  < is total.

PROOF: Given natural numbers m, n, p and q, either m+q < n+p, or m+q=n+p, or n+p < m+q.

**Definition 16.1.22** (Positive). An integer a is positive iff a > 0.

**Theorem 16.1.23.** For any integers a, b and c, we have a < b if and only if a + c < b + c.

- $\langle 1 \rangle 1$ . If a < b then a + c < b + c.
  - $\langle 2 \rangle 1$ . Let: a = [(m, n)], b = [(p, q)] and c = [(r, s)].
  - $\langle 2 \rangle 2$ . Assume: a < b
  - $\langle 2 \rangle 3. \ m+q < n+p$
  - $\langle 2 \rangle 4$ . m + r + q + s < n + r + p + s
  - $\langle 2 \rangle 5. [(m+r, n+s)] < [(p+r, q+s)]$
  - $\langle 2 \rangle 6$ . a+c < b+c

```
\langle 1 \rangle2. If a+c < b+c then a < b.
PROOF: From \langle 1 \rangle1 and Proposition 6.2.6.
```

**Proposition 16.1.24.** Let a, b and c be integers. If 0 < c, then a < b if and only if ac < bc.

```
PROOF:  \langle 1 \rangle 1. \text{ Let: } c = [(r,s)] \\ \langle 1 \rangle 2. \text{ Assume: } 0 < c \\ \langle 1 \rangle 3. s < r \\ \langle 1 \rangle 4. \text{ For all integers } a \text{ and } b, \text{ if } a < b \text{ then } ac < bc \\ \langle 2 \rangle 1. \text{ Let: } a = [(m,n)], b = [(p,q)]. \\ \langle 2 \rangle 2. \text{ Assume: } a < b \\ \langle 2 \rangle 3. m + q < n + p \\ \langle 2 \rangle 4. (m+q)r + (p+n)s < (m+q)s + (p+n)r \\ \text{PROOF: Proposition } 9.4.6, \langle 1 \rangle 3, \langle 2 \rangle 3. \\ \langle 2 \rangle 5. mr + ns + ps + qr < ms + nr + pr + qs \\ \langle 2 \rangle 6. [(mr + ns, ms + nr)] < [(pr + qs, ps + qr)] \\ \langle 2 \rangle 7. ac < bc \\ \langle 1 \rangle 5. \text{ For all integers } a \text{ and } b, \text{ if } ac < bc \text{ then } a < b \\ \text{PROOF: From } \langle 1 \rangle 4 \text{ and Proposition } 6.2.6. \\ \square
```

**Proposition 16.1.25.** Let a be a positive integer. For any integer b, there exists  $k \in \mathbb{N}$  such that b < ak.

```
PROOF: \langle 1 \rangle 1. Case: b \leq 0
PROOF: Take k = 1. \langle 1 \rangle 2. Case: b > 0
PROOF: Take k = b + 1.
```

16.2 The Rationals

**Definition 16.2.1** (Rational Numbers). The set  $\mathbb{Q}$  of rational numbers is the field of fractions over the integers.

**Proposition 16.2.2.** For any integers a and b, we have [(a,1)] = [(b,1)] iff a = b.

Proof: Immediate from definitions.

Henceforth we identify any integer a with the rational number [(a, 1)].

**Proposition 16.2.3.** Addition on the rationals agrees with addition on the integers.

PROOF: 
$$[(a, 1)] + [(b, 1)] = [(a \cdot 1 + b \cdot 1, 1 \cdot 1)] = [(a + b, 1)].$$

**Proposition 16.2.4.** Multiplication on the rationals agrees with multiplication on the integers.

PROOF: 
$$[(a, 1)][(b, 1)] = [(ab, 1)]$$

**Definition 16.2.5.** Define the ordering < on the rationals by: if b and d are positive, then [(a,b)] < [(c,d)] iff ad < bc.

We prove this is well-defined.

#### Proof:

 $\langle 1 \rangle 1$ . For any rational q, there exist integers a, b with b positive such that q = [(a, b)].

PROOF: Since [(a,b)] = [(-a,-b)], and if  $b \neq 0$  then one of b and -b is positive.

 $\langle 1 \rangle$ 2. If b, b', d and d' are positive, [(a,b)] = [(a',b')], and [(c,d)] = [(c',d')], then ad < bc iff a'd' < b'c'.

### Proof:

- $\langle 2 \rangle 1$ . If ad < bc then a'd' < b'c'.
  - $\langle 3 \rangle 1$ . Assume: ad < bc
  - $\langle 3 \rangle 2$ . ab'd < bb'c
  - $\langle 3 \rangle 3$ . a'bd < bb'c
  - $\langle 3 \rangle 4$ . a'd < b'c
  - $\langle 3 \rangle 5$ . a'dd' < b'cd'
  - $\langle 3 \rangle 6$ . a'dd' < b'c'd
  - $\langle 3 \rangle 7$ . a'd' < b'c'
- $\langle 2 \rangle 2$ . If a'd' < b'c' then ad < bc.

PROOF: Similar.

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**Proposition 16.2.6.** The ordering on the rationals agrees with the ordering on the integers.

Proof: We have [(a,1)] < [(b,1)] if and only if a < b.  $\square$ 

**Proposition 16.2.7.** The relation < is a linear ordering on  $\mathbb{Q}$ .

### Proof:

 $\langle 1 \rangle 1$ . < is irreflexive.

PROOF: We never have ab < ab.

- $\langle 1 \rangle 2$ . < is transitive.
  - $\langle 2 \rangle 1$ . Assume: [(a,b)] < [(c,d)] < [(e,f)] where b, d and f are positive.
  - $\langle 2 \rangle 2$ . ad < bc and cf < de
  - $\langle 2 \rangle 3$ . adf < bde

Proof: adf < bcf < bde

- $\langle 2 \rangle 4$ . af < be
- $\langle 1 \rangle 3. < \text{is total.}$

PROOF: For any integers a, b, c, d, we have ad < bc or ad = bc or bc < ad.

**Proposition 16.2.8.** For any rationals r, s and t, we have r < s if and only if r + t < s + t.

Proof:

 $\langle 1 \rangle 1$ . Let: a, b, c, d, e, f be integers with b, d and f positive.

 $\langle 1 \rangle 2$ . [(a,b)] + [(e,f)] < [(c,d)] + [(e,f)] if and only if [(a,b)] < [(c,d)]. PROOF:

$$\begin{split} [(a,b)] + [(e,f)] < [(c,d)] + [(e,f)] &\Leftrightarrow [(af+be,bf)] < [(cf+de,df)] \\ &\Leftrightarrow (af+be)df < (cf+de)bf \\ &\Leftrightarrow afdf+bedf < cfbf+debf \\ &\Leftrightarrow afdf < cfbf \\ &\Leftrightarrow ad < bc \\ &\Leftrightarrow [(a,b)] < [(c,d)] \end{split}$$

**Corollary 16.2.8.1.** For any rational r, we have r < 0 if and only if 0 < -r.

**Definition 16.2.9** (Absolute Value). For any rational r, the absolute value of r is defined by

$$|r| := \begin{cases} -r & \text{if } 0 < -r \\ r & \text{otherwise} \end{cases}$$

**Proposition 16.2.10.** For any rationals r, s and t, if t is positive then r < s iff rt < st.

Proof:

 $\langle 1 \rangle 1$ . Let: r = [(a,b)], s = [(c,d)] and t = [(e,f)] where b, d and f are positive.

 $\langle 1 \rangle 2$ . Assume: 0 < t

 $\langle 1 \rangle 3. \ e > 0$ 

 $\langle 1 \rangle 4$ . rt < st iff r < s

Proof:

$$rt < st \Leftrightarrow [(ae, bf)] < [(ce, df)]$$
  
 $\Leftrightarrow aedf < cebf$   
 $\Leftrightarrow ad < bc$   
 $\Leftrightarrow r < s$ 

Corollary 16.2.10.1. The rationals form an ordered field.

**Proposition 16.2.11.** *Let* p *be a positive rational. For any rational number* r, *there exists*  $k \in \mathbb{N}$  *such that* r < pk.

Proof:

 $\langle 1 \rangle 1$ . Let: p = a/b and r = c/d where a, b and d are positive.

```
\langle 1 \rangle2. PICK k \in \mathbb{N} such that bc < adk PROOF: Proposition 16.1.25. \langle 1 \rangle3. r < pk
```

### Proposition 16.2.12. $\mathbb{Q} \approx \mathbb{N}$

PROOF: Arrange the rationals in order 0/1, 1/1, 1/2, 0/2, -1/2, -1/1, -2/1, -2/2, -2/3, -1/3, 0/3, 1/3, 2/3, etc. then remove all duplicates.  $\Box$ 

### 16.3 The Real Numbers

**Definition 16.3.1** (Cauchy Sequence). A Cauchy sequence is a sequence  $(q_n)$  of rationals such that, for every positive rational  $\epsilon$ , there exists  $k \in \mathbb{N}$  such that  $\forall m, n > k. |q_m - q_n| < \epsilon$ .

**Definition 16.3.2** (Dedekind Cut). A *Dedekind cut* is a set  $x \subseteq \mathbb{Q}$  such that:

- 1.  $\emptyset \neq x \neq \mathbb{Q}$
- 2. x is closed downwards.
- 3. x has no greatest member.

The set  $\mathbb{R}$  of *real numbers* is the set of Dedekind cuts.

**Proposition 16.3.3.** For any rational q, we have  $\{r \in \mathbb{Q} \mid r < q\} \in \mathbb{R}$ .

Proof:

- $\langle 1 \rangle 1$ . Let:  $q \in \mathbb{Q}$
- $\langle 1 \rangle 2$ . Let:  $q \downarrow = \{r \mid r < q\}$
- $\langle 1 \rangle 3. \ q \downarrow \neq \emptyset$

PROOF: We have  $q - 1 \in q \downarrow$ .

 $\langle 1 \rangle 4. \ \ q \downarrow \neq \mathbb{Q}$ 

PROOF: Since  $q \notin q \downarrow$ .

 $\langle 1 \rangle 5$ .  $q \downarrow$  is closed downwards.

PROOF: Trivial.

 $\langle 1 \rangle 6$ .  $q \downarrow$  has no greatest element.

PROOF: For all  $r \in q \downarrow$  we have  $r < (q+r)/2 \in q \downarrow$ .

**Proposition 16.3.4.** For rationals q and r, we have q = r if and only if  $\{s \in \mathbb{Q} \mid s < q\} = \{s \in \mathbb{Q} \mid s < r\}.$ 

### Proof:

- $\langle 1 \rangle 1$ . Let:  $q \downarrow = \{ s \in \mathbb{Q} \mid s < q \}$
- $\langle 1 \rangle 2$ . Let:  $r \downarrow = \{ s \in \mathbb{Q} \mid s < r \}$
- $\langle 1 \rangle 3$ . If q = r then  $q \downarrow = r \downarrow$

PROOF: Trivial.

```
\langle 1 \rangle4. If q < r then q \downarrow \neq r \downarrow
PROOF: We have q \in r \downarrow and q \notin q \downarrow.
\langle 1 \rangle5. If r < q then q \downarrow \neq r \downarrow
PROOF: We have r \in q \downarrow and q \notin q \downarrow.
```

Henceforth we identify a rational q with the real number  $\{r \in \mathbb{Q} \mid r < q\}$ .

**Definition 16.3.5.** Define the ordering < on  $\mathbb{R}$  by: x < y iff  $x \subseteq y$ .

**Proposition 16.3.6.** The ordering on the reals agrees with the ordering on the rationals.

```
Proof:
\langle 1 \rangle 1. Let: q, r \in \mathbb{Q}
\langle 1 \rangle 2. Let: q \downarrow = \{ s \in \mathbb{Q} \mid s < q \}.
\langle 1 \rangle 3. Let: r \downarrow = \{ s \in \mathbb{Q} \mid s < r \}.
            Prove: q < r \text{ iff } q \downarrow \subsetneq r \downarrow
\langle 1 \rangle 4. If q < r then q \downarrow \subseteq r \downarrow
     \langle 2 \rangle 1. Assume: q < r
     \langle 2 \rangle 2. q \downarrow \subseteq r \downarrow
          Proof: If s < q then s < r.
     \langle 2 \rangle 3. \ \ q \downarrow \neq r \downarrow
          Proof: Proposition 16.3.4.
\langle 1 \rangle 5. If q \downarrow \subsetneq r \downarrow then q < r
     \langle 2 \rangle 1. Assume: q \downarrow \subsetneq r \downarrow
     \langle 2 \rangle 2. Pick s \in r \downarrow such that s \notin q \downarrow
     \langle 2 \rangle 3. \ q \leq s < r
```

**Proposition 16.3.7.** The ordering < is a linear ordering on  $\mathbb{R}$ .

```
Proof:
```

```
\langle 1 \rangle 1. < is irreflexive.
```

PROOF: No set is a proper subset of itself.

 $\langle 1 \rangle 2$ . < is transitive.

PROOF: Since the relationship  $\subseteq$  is transitive on the class of all sets.

- $\langle 1 \rangle 3$ . < is total.
  - $\langle 2 \rangle 1$ . Let: x, y be Dedekind cuts.
  - $\langle 2 \rangle 2$ . Assume:  $x \nsubseteq y$ Prove:  $y \subsetneq x$
  - $\langle 2 \rangle 3$ . PICK  $q \in x$  such that  $q \notin y$
  - $\langle 2 \rangle 4$ . Let:  $r \in y$ Prove:  $r \in x$
  - $\langle 2 \rangle 5. \ q \not\leq r$

PROOF: Since y is closed downwards.

- $\langle 2 \rangle 6$ . r < q
- $\langle 2 \rangle 7. \ r \in x$

PROOF: Since x is closed downwards.

**Proposition 16.3.8.** Any bounded nonempty subset of  $\mathbb{R}$  has a least upper bound.

### Proof:

- $\langle 1 \rangle 1$ . Let: A be a bounded nonempty subset of  $\mathbb{R}$ .
- $\langle 1 \rangle 2$ .  $\bigcup A$  is a Dedekind cut.
  - $\langle 2 \rangle 1. \bigcup A \neq \emptyset$ 
    - $\langle 3 \rangle 1$ . Pick $x \in A$
    - $\langle 3 \rangle 2$ . Pick  $q \in x$
    - $\langle 3 \rangle 3. \ q \in \bigcup A$
  - $\langle 2 \rangle 2$ .  $\bigcup A \neq \mathbb{Q}$ 
    - $\langle 3 \rangle 1$ . PICK an upper bound u for A
    - $\langle 3 \rangle 2$ . Pick  $q \notin u$ Prove:  $q \notin \bigcup A$
    - $\langle 3 \rangle 3$ . Assume: for a contradiction  $q \in \bigcup A$
    - $\langle 3 \rangle 4$ . PICK  $x \in A$  such that  $q \in x$
    - $\langle 3 \rangle 5. \ x \leq u$
    - $\langle 3 \rangle 6. \ q \in u$
    - $\langle 3 \rangle$ 7. Q.E.D.

PROOF: This is a contradiction.

- $\langle 2 \rangle 3$ .  $\bigcup A$  is closed downwards.
  - $\langle 3 \rangle 1$ . Let:  $q \in \bigcup A$  and r < q
  - $\langle 3 \rangle 2$ . PICK  $x \in A$  such that  $q \in x$
  - $\langle 3 \rangle 3. \ r \in x$
  - $\langle 3 \rangle 4. \ r \in \bigcup A$
- $\langle 2 \rangle 4$ .  $\bigcup A$  has no greatest element.
  - $\langle 3 \rangle 1$ . Let:  $q \in \bigcup A$
  - $\langle 3 \rangle 2$ . PICK  $x \in A$  such that  $q \in x$
  - $\langle 3 \rangle 3$ . Pick  $r \in x$  such that q < r
  - $\langle 3 \rangle 4. \ r \in \bigcup A$
- $\langle 1 \rangle 3$ .  $\bigcup A$  is an upper bound for A.

PROOF: For all  $x \in A$  we have  $x \subseteq \bigcup A$ .

 $\langle 1 \rangle 4$ . For any upper bound u for  $\bigcup A$  we have  $\bigcup A \leq u$ .

PROOF: If  $\forall x \in A.x \subseteq u$  we have  $\bigcup A \subseteq u$ .

**Definition 16.3.9** (Addition). Define addition + on the reals by

$$x + y := \{q + r \mid q \in x, r \in y\}$$
.

We prove this is well-defined.

### Proof:

 $\langle 1 \rangle 1$ . Let:  $x, y \in \mathbb{R}$ 

PROVE: X + y is a Dedekind cut.

```
\langle 1 \rangle 2. \ x + y \neq \emptyset
   PROOF: Pick q \in x and r \in y; then q + r \in x + y.
\langle 1 \rangle 3. \ x + y \neq \mathbb{Q}
   \langle 2 \rangle 1. PICK q \notin x and r \notin y
           PROVE: q + r \notin x + y
   \langle 2 \rangle 2. Assume: for a contradiction q + r \in x + y
   \langle 2 \rangle 3. Pick q' \in x and r' \in y such that q + r = q' + r'
   \langle 2 \rangle 4. q' < q and r' < r
   \langle 2 \rangle 5. q' + r' < q + r
   \langle 2 \rangle 6. Q.E.D.
      PROOF: This is a contradiction.
\langle 1 \rangle 4. x + y is closed downwards.
   \langle 2 \rangle 1. Let: q \in x and r \in y
   \langle 2 \rangle 2. Let: s < q + r
            PROVE: s \in x + y
   \langle 2 \rangle 3. \ s - r < q
   \langle 2 \rangle 4. \ s - r \in x
   \langle 2 \rangle 5. s = (s - r) + r \in x + y
\langle 1 \rangle 5. x + y has no greatest element.
   \langle 2 \rangle 1. Let: q \in x and r \in y
            PROVE: There exists s \in x + y such that q + r < s
   \langle 2 \rangle 2. Pick q' \in x and r' \in y such that q < q' and r < r'
   \langle 2 \rangle 3. \ \ q + r < q' + r' \in x + y
```

**Proposition 16.3.10.** Addition on the reals agrees with addition on the rationals.

```
\begin{split} &\langle 1 \rangle 1. \text{ Let: } q,r \in \mathbb{Q} \\ &\langle 1 \rangle 2. \ q \downarrow + r \downarrow \subseteq (q+r) \downarrow \\ &\text{Proof: If } s_1 < q \text{ and } s_2 < r \text{ then } s_1 + s_2 < q + r. \\ &\langle 1 \rangle 3. \ (q+r) \downarrow \subseteq q \downarrow + r \downarrow \\ &\langle 2 \rangle 1. \text{ Let: } s < q + r \\ &\langle 2 \rangle 2. \ s - r < q \\ &\langle 2 \rangle 3. \text{ Pick } t \text{ such that } s - r < t < q \end{split}
```

 $\langle 2 \rangle 4$ . s - t < r

 $\langle 2 \rangle 5. \ \ s = t + (s - t) \in q \downarrow +r \downarrow$ 

Proposition 16.3.11. Addition is associative.

Proof:

Proof:

$$x + (y + z) = \{q + r \mid q \in x, r \in y + z\}$$

$$= \{q + s_1 + s_2 \mid q \in x, s_1 \in y, s_2 \in z\}$$

$$= \{r + s_2 \mid r \in x + y, s_2 \in z\}$$

$$= (x + y) + z$$

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Proposition 16.3.12. Addition is commutative.

Proof:

$$x + y = \{q + r \mid q \in x, r \in y\}$$

$$= \{r + q \mid r \in y, q \in x\}$$

$$= y + x$$

**Proposition 16.3.13.** For any  $x \in \mathbb{R}$  we have x + 0 = x.

Proof:

 $\langle 1 \rangle 1$ .  $x + 0 \subseteq x$ 

PROOF: If  $q \in x$  and r < 0 then q + r < q so  $q + r \in x$ .

- $\langle 1 \rangle 2. \ x \subseteq x + 0$ 
  - $\langle 2 \rangle 1$ . Let:  $q \in x$
  - $\langle 2 \rangle 2$ . Pick  $r \in x$  such that q < r.

PROOF: x has no greatest element.

- $\langle 2 \rangle 3. \ q-r < 0$
- $\langle 2 \rangle 4. \ q = r + (q r) \in x + 0$

**Definition 16.3.14.** For  $x \in \mathbb{R}$ , define  $-x := \{q \in \mathbb{Q} \mid \exists r > q. -r \notin x\}$ .

**Proposition 16.3.15.** For all  $x \in \mathbb{R}$  we have  $-x \in \mathbb{R}$ .

Proof:

- $\langle 1 \rangle 1$ . Let:  $x \in \mathbb{R}$
- $\langle 1 \rangle 2. -x \neq \emptyset$ 
  - $\langle 2 \rangle 1$ . Pick  $s \notin x$
  - $\langle 2 \rangle 2$ .  $-s-1 \in -x$
- $\langle 1 \rangle 3. -x \neq \mathbb{Q}$ 
  - $\langle 2 \rangle 1$ . Pick  $s \in x$

Prove:  $-s \notin -x$ 

- $\langle 2 \rangle 2$ . Assume: for a contradiction  $-s \in -x$
- $\langle 2 \rangle 3$ . PICK r > -s such that  $-r \notin x$
- $\langle 2 \rangle 4$ . -r < s
- $\langle 2 \rangle 5$ . Q.E.D.

PROOF: This contradicts the fact that x is closed downwards.

 $\langle 1 \rangle 4$ . -x is closed downwards.

PROOF: Immediate from definition.

- $\langle 1 \rangle 5$ . -x has no greatest element.
  - $\langle 2 \rangle 1$ . Let:  $q \in -x$
  - $\langle 2 \rangle 2$ . Pick r > q such that  $-r \notin x$
  - $\langle 2 \rangle 3$ . Pick s such that q < s < r
- $\langle 2 \rangle 4. \ s \in -x$

**Lemma 16.3.16.** Let p be a positive rational number. For any real number x, there exists a rational  $q \in x$  such that  $p + q \notin x$ .

```
Proof:
\langle 1 \rangle 1. Pick q_0 \in x
\langle 1 \rangle 2. There exists k \in \mathbb{N} such that q_0 + kp \notin x
   \langle 2 \rangle 1. Pick q_1 \notin x
   \langle 2 \rangle 2. PICK k \in \mathbb{N} such that q_1 - q_0 < pk
       Proof: Proposition 16.2.11.
    \langle 2 \rangle 3. \ q_1 < q_0 + kp
   \langle 2 \rangle 4. \ q_0 + kp \notin x
\langle 1 \rangle 3. Let: k be the least natural number such that q_0 + kp \notin x
\langle 1 \rangle 4. \ k \neq 0
   Proof: \langle 1 \rangle 1
\langle 1 \rangle 5. Let: q = q_0 + (k-1)p
\langle 1 \rangle 6. q \in x and q + p \notin x.
Proposition 16.3.17. For every real x we have x + (-x) = 0.
Proof:
\langle 1 \rangle 1. Let: x be a real number.
\langle 1 \rangle 2. x + (-x) \subseteq 0
   \langle 2 \rangle 1. Let: q_1 \in x and q_2 \in -x
   \langle 2 \rangle 2. PICK r > q_2 such that -r \notin x
   \langle 2 \rangle 3. \ q_1 < -r
   \langle 2 \rangle 4. r < -q_1
   \langle 2 \rangle 5. q_2 < -q_1
   \langle 2 \rangle 6. \ q_1 + q_2 < 0
\langle 1 \rangle 3. \ 0 \subseteq x + (-x)
   \langle 2 \rangle 1. Let: p < 0
   \langle 2 \rangle 2. 0 < -p
   \langle 2 \rangle 3. Pick q \in x such that q - p/2 \notin x
       Proof: Lemma 16.3.16.
    \langle 2 \rangle 4. Let: s = p/2 - q
   \langle 2 \rangle 5. -s \notin x
   \langle 2 \rangle 6. \ p - q < s
   \langle 2 \rangle 7. \ p-q \in -x
    \langle 2 \rangle 8. \ p \in x + (-x)
```

Corollary 16.3.17.1. The reals form an Abelian group under addition.

**Proposition 16.3.18.** For any reals x, y and z, we have x < y if and only if x + z < y + z.

```
\begin{array}{l} \langle 1 \rangle 1. \ \forall x,y,z \in \mathbb{R}. x \leq y \Rightarrow x+z \leq y+z \\ \langle 2 \rangle 1. \ \text{Let:} \ x,y,z \in \mathbb{R} \\ \langle 2 \rangle 2. \ \text{Assume:} \ x \leq y \\ \langle 2 \rangle 3. \ \text{For all} \ q \in x \ \text{and} \ r \in z \ \text{we have} \ q+r \in y+z \end{array}
```

 $\langle 1 \rangle 2$ .  $\forall x, y, z \in \mathbb{R}.x + z = y + z \Leftrightarrow x = y$ PROOF: Proposition 13.1.4.

 $\langle 1 \rangle 3. \ \forall x, y, z \in \mathbb{R}. x < y \Rightarrow x + z < y + z$ 

 $\langle 1 \rangle 4$ . Q.E.D.

Proof: Proposition 6.2.6.

П

**Definition 16.3.19** (Absolute Value). The *absolute value* of a real number x is defined to be

$$|x| = \begin{cases} x & \text{if } 0 \le x \\ -x & \text{if } x < 0 \end{cases}$$

**Definition 16.3.20** (Multiplication). Define *multiplication*  $\cdot$  on  $\mathbb{R}$  as follows:

• If x and y are non-negative then

$$xy = 0 \cup \{rs \mid 0 \le r \in x \land 0 \le s \in y\} .$$

• If x and y are both negative then

$$xy = (-x)(-y) .$$

• If one of x and y is negative and one is non-negative then

$$xy = -(|x||y|) .$$

We prove this is well-defined.

Proof:

 $\langle 1 \rangle 1$ . Let: x and y be non-negative reals.

PROVE: xy is real.

 $\langle 1 \rangle 2. \ xy \neq \emptyset$ 

PROOF: Since  $-1 \in xy$ .

 $\langle 1 \rangle 3. \ xy \neq \mathbb{Q}$ 

 $\langle 2 \rangle 1$ . Pick  $r \notin x$  and  $s \notin y$ 

Prove:  $rs \notin xy$ 

 $\langle 2 \rangle 2$ .  $0 \le r$  and  $0 \le s$ 

PROOF: Since  $0 \subseteq x$  and  $0 \subseteq y$ .

- $\langle 2 \rangle 3$ . Assume: for a contradiction  $rs \in xy$
- $\langle 2 \rangle 4$ . Pick r' and s' such that  $0 \leq r' \in x$ ,  $0 \leq s' \in y$  and rs = r's'
- $\langle 2 \rangle 5. \ r' < r$
- $\langle 2 \rangle 6. \ s' < s$
- $\langle 2 \rangle 7$ . r's' < rs
- $\langle 2 \rangle 8$ . Q.E.D.

PROOF: This is a contradiction.

- $\langle 1 \rangle 4$ . xy is closed downwards.
  - $\langle 2 \rangle 1$ . Let:  $q \in xy$  and r < q

```
\langle 2 \rangle 2. Case: q \in 0
      PROOF: Then r < q < 0 so r \in xy
   \langle 2 \rangle 3. Case: q = s_1 s_2 where 0 \le s_1 \in x and 0 \le s_2 \in y
      \langle 3 \rangle 1. Assume: w.l.o.g. 0 \le r
      \langle 3 \rangle 2. 0 < s_1 and 0 < s_2
      \langle 3 \rangle 3. \ r/s_2 < s_1
      \langle 3 \rangle 4. \ r/s_2 \in x
      \langle 3 \rangle 5. r = (r/s_2)s_2 \in xy
\langle 1 \rangle 5. xy has no greatest element.
   \langle 2 \rangle 1. Let: q \in xy
   \langle 2 \rangle 2. Case: q \in 0
      Proof: q < q/2 \in 0
   \langle 2 \rangle 3. Case: q = rs where 0 \le r \in x and 0 \le s \in y
      \langle 3 \rangle 1. Pick r' and s' with r < r' \in x and s < s' \in y
      \langle 3 \rangle 2. q < r's' \in xy
П
Proposition 16.3.21. Multiplication is commutative.
PROOF: Immediate from definition.
Proposition 16.3.22. Multiplication is associative.
Proof:
\langle 1 \rangle 1. For non-negative reals x, y and z, we have x(yz) = (xy)z
   PROOF: It computes to 0 \cup \{qrs \mid 0 \le q \in x, 0 \le r \in y, 0 \le s \in z\}.
\langle 1 \rangle 2. For all reals x, y and z, we have x(yz) = (xy)z
   PROOF: It is equal to |x||y||z| if an even number of them are negative, and
   -(|x||y||z|) otherwise.
Proposition 16.3.23. Multiplication distributes over addition.
\langle 1 \rangle 1. For all non-negative reals x, y and z, we have x(y+z) = xy + xz
   \langle 2 \rangle 1. Let: x, y and z be non-negative reals.
   \langle 2 \rangle 2. x(y+z) \subseteq xy+xz
      \langle 3 \rangle 1. Let: q \in x(y+z)
      \langle 3 \rangle 2. Case: q < 0
         PROOF: Then we have q/2 \in xy and q/2 \in xz so q \in xy + xz.
      \langle 3 \rangle 3. Case: q = rs where 0 \le r \in x and 0 \le s \in y + z
         \langle 4 \rangle 1. PICK s_1 \in y and s_2 \in z such that s = s_1 + s_2
         \langle 4 \rangle 2. \ rs_1 \in xy
            PROOF: If s_1 < 0 then rs_1 < 0 so rs_1 \in xy. If 0 \le s_1 then we also
            have rs_1 \in xy.
         \langle 4 \rangle 3. \ rs_2 \in xz
            PROOF: Similar.
         \langle 4 \rangle 4. \ \ q \in xy + xz
```

PROOF: Since  $q = rs_1 + rs_2$ .

- $\langle 2 \rangle 3. \ xy + xz \subseteq x(y+z)$ 
  - $\langle 3 \rangle 1$ . Let:  $q \in xy$  and  $r \in xz$ .

PROVE:  $q + r \in x(y + z)$ 

 $\langle 3 \rangle 2$ . Case: q < 0 and r < 0

PROOF: Then q + r < 0 so  $q + r \in x(y + z)$ .

- $\langle 3 \rangle 3$ . Case: q < 0 and  $r = r_1 r_2$  where  $0 \le r_1 \in x$  and  $0 \le r_2 \in z$ 
  - $\langle 4 \rangle 1. \ q + r < r$
  - $\langle 4 \rangle 2. \ q + r \in xz$
  - $\langle 4 \rangle 3$ . Assume: w.l.o.g.  $0 \le q + r$

PROOF: Otherwise  $q + r \in x(y + z)$  immediately.

- $\langle 4 \rangle 4$ . PICK  $s_1, s_2$  with  $0 \leq s_1 \in x$ ,  $0 \leq s_2 \in y$  and  $q + r = s_1 s_2$
- $\langle 4 \rangle 5. \ s_2 \in y + z$

PROOF: Since  $0 \in z$  so  $s_2 = s_2 + 0 \in y + z$ .

- $\langle 4 \rangle 6. \ q+r \in x(y+z)$
- $\langle 3 \rangle 4$ . Case:  $q = q_1 q_2$  where  $0 \le q_1 \in x$  and  $0 \le q_2 \in y$  and r < 0 Proof: Similar.
- $\langle 3 \rangle$ 5. Case:  $q = q_1q_2$  where  $0 \le q_1 \in x$  and  $0 \le q_2 \in y$  and  $r = r_1r_2$  where  $0 \le r_1 \in x$  and  $0 \le r_2 \in z$ 
  - $\langle 4 \rangle 1$ . Assume: w.l.o.g.  $q_1 \leq r_1$
  - $\langle 4 \rangle 2. \ \ q + r \le r_1(q_2 + r_2) \in x(y + z)$
- $\langle 1 \rangle$ 2. For any negative real x and non-negative reals y and z, we have x(y+z)=xy+xz

Proof:

$$x(y+z) = -(-x)(y+z) = -((-x)y + (-x)z)$$

$$= -((-x)y) - ((-x)z)$$

$$= xy + xz$$
(\langle 1\rangle 1)

- $\langle 1 \rangle$ 3. For any non-negative real x and reals y and z with one negative and one non-negative, we have x(y+z)=xy+xz
  - $\langle 2 \rangle 1$ . Assume: w.l.o.g. y is negative and z is non-negative.
  - $\langle 2 \rangle 2$ . Case:  $0 \le y + z$

$$xy + xz = xy + x(-y + y + z)$$

$$= -(x(-y)) + x(-y + y + z)$$

$$= -(x(-y)) + x(-y) + x(y + z)$$

$$= x(y + z)$$
(\langle 1\rangle 1)

- $\langle 2 \rangle 3$ . Case: y + z < 0
  - $\langle 3 \rangle 1. -y z > 0$
  - $\langle 3 \rangle 2$ . -y = z y z
  - $\langle 3 \rangle 3$ . xy + xz = x(y+z)

Proof:

$$xy + xz = -(x(-y)) + xz$$

$$= -(x(z - y - z)) + xz$$

$$= -(xz + x(-y - z)) + xz \qquad (\langle 1 \rangle 1)$$

$$= -xy - x(-y - z) + xz$$

$$= -x(-y - z)$$

$$= x(y + z)$$

 $\langle 1 \rangle 4$ . For any non-negative real x and negative reals y and z, we have x(y+z)=xy+xz

Proof:

$$x(y+z) = -x(-y-z)$$

$$= -(x(-y) + x(-z))$$

$$= -x(-y) - x(-z)$$

$$= xy + xz$$

$$(\langle 1 \rangle 1)$$

- $\langle 1 \rangle$ 5. For any negative real x and reals y and z with one negative and one non-negative, we have x(y+z)=xy+xz
  - $\langle 2 \rangle 1$ . Assume: w.l.o.g. y is negative and z is non-negative.
  - $\langle 2 \rangle 2$ . Case:  $0 \le y + z$

Proof:

$$x(y+z) = -((-x)(y+z))$$

$$= -((-x)y + (-x)z)$$

$$= -((-x)y) - ((-x)z)$$

$$= (-x)(-y) - ((-x)z)$$

$$= xy + xz$$
(\langle 1\rangle 3)

 $\langle 2 \rangle 3$ . Case: y + z < 0

Proof:

$$x(y+z) = (-x)(-y-z) = (-x)(-y) + (-x)(-z) = xy + xz$$
 (\langle 1\rangle 3)

 $\langle 1 \rangle$ 6. For any negative reals x, y and z, we have x(y+z) = xy + xz Proof:

$$x(y+z) = (-x)(-y-z)$$

$$= (-x)(-y) + (-x)(-z)$$

$$= xy + xz$$

$$(\langle 1 \rangle 1)$$

**Proposition 16.3.24.** For any real x we have x1 = x.

- $\langle 1 \rangle 1$ . Case:  $0 \le x$ 
  - $\langle 2 \rangle 1. \ x1 \subseteq x$ 
    - $\langle 3 \rangle 1$ . Let:  $q \in x1$

$$\langle 3 \rangle 2. \text{ CASE: } q < 0$$

$$\text{PROOF: Then } q \in x \text{ because } 0 \leq x.$$

$$\langle 3 \rangle 3. \ q = rs \text{ where } 0 \leq r \in x \text{ and } 0 \leq s < 1$$

$$\text{PROOF: Then } q < r \text{ so } q \in x.$$

$$\langle 2 \rangle 2. \ x \subseteq x1$$

$$\langle 3 \rangle 1. \ \text{Let: } q \in x$$

$$\langle 3 \rangle 2. \ \text{Assume: w.l.o.g. } 0 \leq q$$

$$\langle 3 \rangle 3. \ \text{PICK } r \text{ such that } q < r \in x$$

$$\langle 3 \rangle 4. \ 0 \leq q/r < 1$$

$$\langle 3 \rangle 5. \ q = r(q/r) \in x1$$

$$\langle 1 \rangle 2. \ \text{Case: } x < 0$$

$$\text{PROOF: } x1 = -((-x)1)$$

$$= -(-x)$$

$$= x$$

$$\langle (1) \rangle 1$$

**Lemma 16.3.25.** Let  $x \in \mathbb{R}$  and c be a positive rational. Then there exists  $a \in x$  and a non-least rational upper bound b for x such that b - a = c.

#### PROOF:

- (1)1. PICK  $a_1 \in x$  such that if x has a rational supremum s then  $a_1 > s c$
- $\langle 1 \rangle 2$ . There exists a natural number n such that  $a_1 + nc$  is an upper bound for x.
  - $\langle 2 \rangle 1$ . PICK a non-least upper bound  $b_1$  for x.
  - $\langle 2 \rangle 2$ . PICK a natural number n such that  $nc > b_1 a_1$

Proof: Proposition 16.2.11.

- $\langle 2 \rangle 3$ .  $a_1 + nc > b_1$
- $\langle 2 \rangle 4$ .  $a_1 + nc$  is an upper bound for x.
- $\langle 1 \rangle$ 3. Let: k be the least natural number such that  $a_1 + kc$  is an upper bound for x.
- $\langle 1 \rangle 4. \ a_1 + (k-1)c \in x$
- $\langle 1 \rangle 5$ .  $a_1 + kc$  is not the supremum of x.
  - $\langle 2 \rangle 1$ . Assume: for a contradiction  $a_1 + kc$  is the supremum of x.
  - $\langle 2 \rangle 2$ .  $a_1 > a_1 + (k-1)c$

Proof:  $\langle 1 \rangle 1$ 

 $\langle 2 \rangle 3$ . Q.E.D.

PROOF: This is a contradiction.

- $\langle 1 \rangle 6$ . Let:  $a = a_1 + (k-1)c$
- $\langle 1 \rangle 7$ . Let:  $b = a_1 + kc$
- $\langle 1 \rangle 8. \ b-a=c$

 $\Box$ 

**Proposition 16.3.26.** For any non-zero real x, there exists a real y such that xy = 1.

```
\langle 1 \rangle 1. Case: 0 < x
    \langle 2 \rangle 1. Let: y = \{q \in \mathbb{Q} \mid q \leq 0\} \cup \{u^{-1} \mid u \text{ is an upper bound for } x \text{ but not the supremum of } x\}
   \langle 2 \rangle 2. y is a real number.
       \langle 3 \rangle 1. \ y \neq \emptyset
           PROOF: Since 0 \in y.
       \langle 3 \rangle 2. \ y \neq \mathbb{Q}
           \langle 4 \rangle 1. PICK q \in x such that 0 < q
           \langle 4 \rangle 2. \ q^{-1} \notin y
       \langle 3 \rangle 3. y is closed downwards.
           \langle 4 \rangle 1. Let: q \in y and r < q
                    Prove: r \in y
           \langle 4 \rangle 2. Assume: w.l.o.g. 0 < r
           \langle 4 \rangle 3. q^{-1} is a non-least upper bound for x. \langle 4 \rangle 4. q^{-1} < r^{-1}
           \langle 4 \rangle 5. r^{-1} is a non-least upper bound for x.
           \langle 4 \rangle 6. \ r \in y
       \langle 3 \rangle 4. y has no greatest element.
           \langle 4 \rangle 1. Let: q \in y
                    PROVE: There exists r \in y such that q < r
           \langle 4 \rangle 2. Case: q \leq 0
               \langle 5 \rangle 1. PICK a non-least upper bound u for x.
               \langle 5 \rangle 2. \ q < u^{-1} \in x
           \langle 4 \rangle 3. Case: q = u^{-1} where u is a non-least upper bound for x.
               \langle 5 \rangle1. PICK a non-least upper bound v with v < u
               \langle 5 \rangle 2. \ u^{-1} < v^{-1} \in y
    \langle 2 \rangle 3. \ 0 < y
    \langle 2 \rangle 4. xy \subseteq 1
       \langle 3 \rangle 1. Let: q \in xy
       \langle 3 \rangle 2. Assume: w.l.o.g. 0 < q
       \langle 3 \rangle 3. Pick 0 < r \in x and 0 < s \in y such that q = rs
       \langle 3 \rangle 4. \ s^{-1} is a non-least upper bound for x
       \langle 3 \rangle 5. \ r < s^{-1}
       \langle 3 \rangle 6. \ rs < 1
    \langle 2 \rangle 5. 1 \subseteq xy
       \langle 3 \rangle 1. Let: q < 1
                Prove: q \in xy
       \langle 3 \rangle 2. Assume: w.l.o.g. 0 < q
       \langle 3 \rangle 3. PICK a_1 with 0 < a_1 \in x
       \langle 3 \rangle 4. \ (1-q)a_1 > 0
       \langle 3 \rangle 5. Pick a \in x and a non-least upper bound w of x such that w - a =
                (1-q)a_1
           Proof: Lemma 16.3.25.
       \langle 3 \rangle 6. \ \ w - a < (1 - q)w
       \langle 3 \rangle 7. qw < a
       \langle 3 \rangle 8. \ w < a/q
       \langle 3 \rangle 9. a/q is a non-least upper bound for x
```

 $\langle 3 \rangle 10. \ q/a \in y$   $\langle 3 \rangle 11. \ q \in xy$   $\langle 1 \rangle 2. \ \text{Case:} \ x < 0$   $\langle 2 \rangle 1. \ \text{Pick} \ y \ \text{such that} \ (-x)y = 1$   $\text{Proof:} \ \langle 1 \rangle 1$   $\langle 2 \rangle 2. \ x(-y) = 1$ 

**Proposition 16.3.27.** For real numbers x, y and z, if 0 < z then x < y if and only if xz < yz.

#### Proof:

- $\langle 1 \rangle 1$ . For any real numbers x, y and z, if 0 < z and x < y then xz < yz
  - $\langle 2 \rangle 1$ . Let: x, y and z be real numbers.
  - $\langle 2 \rangle 2$ . Assume: 0 < z and x < y.
  - $\langle 2 \rangle 3. \ y = x + (y x)$
  - $\langle 2 \rangle 4$ . y x > 0
  - $\langle 2 \rangle 5$ . (y-x)z > 0
  - $\langle 2 \rangle 6$ . yz > xz

Proof:

$$yz = (x + (y - x))z$$
$$= xz + (y - x)z$$
$$> xz$$

 $\langle 1 \rangle 2$ . For any real numbers x, y and z, if 0 < z and xz < yz then x < y PROOF: Proposition 6.2.6.

Corollary 16.3.27.1. The real numbers form a complete ordered field.

#### Proposition 16.3.28.

$$(0,1) \approx \mathbb{R}$$

PROOF: The function  $f(x) = (2x-1)/(x-x^2)$  is a bijection between (0,1) and  $\mathbb{R}$ .  $\square$ 

#### Proposition 16.3.29.

$$|\mathbb{R}| = 2^{\aleph_0}$$

Proof:

 $\langle 1 \rangle 1. \ (0,1) \leqslant 2^{\mathbb{N}}$ 

PROOF: The function H where H(x)(n) is the nth binary digit of the binary expansion of x is an injection.

 $\langle 1 \rangle 2. \ 2^{\mathbb{N}} \preccurlyeq \mathbb{R}$ 

PROOF: Map f to the real number in [0,1/9] whose n+1st decimal digit is f(n).

**Proposition 16.3.30.** The set of algebraic numbers is countable.

Proof:	: There	are c	countab	oly ma	ny i	nteger	polyn	omials,	each	with	finitely	many
roots.												

Corollary 16.3.30.1. There are uncountably many transcendental numbers.

**Proposition 16.3.31.** Let A be a set of disks in the plane, no two of which intersect. Then A is countable.

PROOF: Every circle includes a point with rational coordinates. Define  $f:\{q\in\mathbb{Q}^2\mid\exists C\in A.q\in C\}\rightarrow A$  by f(q)=C iff  $q\in C$ . Then f is surjective.  $\square$ 

**Proposition 16.3.32.** There exists an uncountable set of circles in the plane that do not intersect.

Proof: The set of all circles with origin O is uncountable.  $\square$ 

## Chapter 17

## Real Analysis

**Theorem 17.0.1** (Weierstrass). Let  $a, b \in \mathbb{R}$  with a < b. Let  $f : [a, b] \to \mathbb{R}$  be continuous. For every  $\epsilon > 0$ , there exists a polynomial p such that  $\forall x \in [a, b].|f(x) - p(x)| < \epsilon$ .

**Theorem 17.0.2** (Bolzano-Weierstrass). Every bounded sequence in  $\mathbb{R}^n$  has a convergent subsequence.

## 17.1 Step Functions

**Definition 17.1.1** (Step Function). A *step function* on  $\mathbb{R}$  is a function  $f: \mathbb{R} \to \mathbb{R}$ , where  $[a_1, b_1), \ldots, [a_n, b_n)$  are disjoint half-open intervals, such that f is constant on each  $[a_i, b_i)$ , and zero outside them.

**Definition 17.1.2** (Basic Representation). Let f be a step function. The *basic representation* of f is defined as follows.

Let  $a_0, a_1, \ldots, a_n$  be the points of discontinuity of f. For  $k = 1, \ldots, n$ , let  $\alpha_k = f(a_{k-1})$  and  $g_k$  be the characteristic function of  $[a_{k-1}, a_k)$ . Then the basic representation of f is

$$f = \alpha_1 g_1 + \dots + \alpha_n g_n .$$

**Proposition 17.1.3.** If f and g are step functions then  $\lambda x.f(x) + g(x)$  is a step function.

**Proposition 17.1.4.** *If*  $c \in \mathbb{R}$  *and* f *is a step function then*  $\lambda x.cf(x)$  *is a step function.* 

**Proposition 17.1.5.** If f is a step function then  $\lambda x.|f(x)|$  is a step function.

**Proposition 17.1.6.** If f and g are step functions then  $\lambda x. \min(f(x), g(x))$  and  $\lambda x. \max(f(x), g(x))$  are step functions.

**Proposition 17.1.7.** If f is a step function and  $c \in \mathbb{R}$  then  $\lambda x. f(x-c)$  is a step function.

**Definition 17.1.8** (Support). Given  $f: \mathbb{R} \to \mathbb{R}$ , the *support* of f is

$$\operatorname{supp} f := \{ x \in \mathbb{R} \mid f(x) \neq 0 \} .$$

**Definition 17.1.9** (Integral of a Step Function). Given a step function f, define  $\int f = \int f(x)dx \in \mathbb{R}$  as follows.

Let  $f(x) = \lambda_1 f_1(x) + \cdots + \lambda_n f_n(x)$ , where  $f_i$  is the characteristic function of  $[a_i, b_i)$ . Then

$$\int f = \lambda_1(b_1 - a_1) + \dots + \lambda_n(b_n - a_n) .$$

We prove this is well defined.

Proof:

- $\langle 1 \rangle 1$ . Let:  $f = \lambda_1 f_1(x) + \dots + \lambda_m f_m(x) = \mu_1 g_1(x) + \dots + \mu_n g_n(x)$ , where  $f_i$  is the characteristic function of  $[a_i, b_i)$  and  $g_i$  is the characteristic function of  $[c_i, d_i)$ , with  $a_1 < b_1 \le a_2 < b_2 \le \dots \le a_m < b_m$  and  $c_1 < d_1 \le c_2 < d_2 \le \dots \le c_n < d_n$ .
- $\langle 1 \rangle 2$ . Assume: w.l.o.g. none of the  $\lambda_i$  or  $\mu_i$  is zero.
- $\langle 1 \rangle 3$ . Assume: w.l.o.g. we never have  $\lambda_i = \lambda_{i+1}$  and  $b_i = a_{i+1}$ , and we never have  $\mu_i = \mu_{i+1}$  and  $d_i = c_{i+1}$ .
- $\langle 1 \rangle 4$ . We have m = n and for all i,  $a_i = b_i$  and  $c_i = d_i$  and  $\lambda_i = \mu_i$ .
  - $\langle 2 \rangle 1$ . Assume: w.l.o.g.  $m \leq n$
  - $\langle 2 \rangle 2$ . Assume: as induction hypothesis  $a_i = b_i$ ,  $c_i = d_i$ ,  $\lambda_i = \mu_i$  for  $i = 1, \ldots, k$  with k < m.
  - $\langle 2 \rangle 3. \ a_{k+1} = b_{k+1}$

PROOF: It is  $\inf\{x \in (a_k, +\infty) : f(x) \neq 0\}$  (or  $\inf\{x \in \mathbb{R} : f(x) \neq 0\}$  if k = 0).

 $\langle 2 \rangle 4. \ \lambda_{k+1} = \mu_{k+1}$ 

PROOF: It is  $f(a_{k+1})$ .

 $\langle 2 \rangle 5.$   $c_{k+1} = d_{k+1}$ 

PROOF: It is  $\sup\{x \in (a_{k+1}, +\infty) : f(x) = \lambda_{k+1}\}.$ 

 $\langle 2 \rangle 6. \ m = n$ 

PROOF: For all  $x > b_m$  we have f(x) = 0.

**Proposition 17.1.10.** If f and g are step functions then  $\int (f(x) + g(x))dx = \int f + \int g$ .

**Proposition 17.1.11.** If f is a step function then  $\int cf(x)dx = c \int f$ .

**Proposition 17.1.12.** If f and g are step functions and  $\forall x. f(x) \leq g(x)$  then  $\int f \leq \int g$ .

PROOF: We have  $g(x) - f(x) \ge 0$  for all x and so  $\int (g(x) - f(x)) dx \ge 0$ .  $\square$ 

**Proposition 17.1.13.** If f is a step function then  $\left| \int f \right| \leq \int |f(x)| dx$ .

**Proposition 17.1.14.** *If* f *is a step function and*  $c \in \mathbb{R}$  *then*  $\int f(x-c)dx = \int f$ .

**Lemma 17.1.15.** Let f be a step function with supp  $f \subseteq [a_1, b_1) \cup \cdots \cup [a_n, b_n)$ . Let M be a constant. If  $\forall x. |f(x)| < M$  then

$$\int |f(x)|dx \le M \sum_{k=1}^{n} (b_k - a_k) .$$

**Lemma 17.1.16.** Let  $\{[a_i,b_i) \mid i \in \mathbb{N}\}$  be a partition of [a,b). Then

$$\sum_{i=0}^{\infty} (b_i - a_i) = b - a .$$

Proof:

 $\langle 1 \rangle 1$ . For all  $c \in (a, b]$  we have  $\{[a_i, b_i) \cap [a, c) \mid i \in \mathbb{N}\}$  is a partition of [a, c).

$$\langle 1 \rangle 2$$
. For  $c \in (a, b]$  and  $n \in \mathbb{N}$ ,

LET: 
$$b_{cn} := min(b_n, c)$$
  
(1)3. LET:  $S = \{c \in (a, b] \mid \sum_{a_n < b_{cn}} (b_{cn} - a_n) = c - a\}$ 

 $\langle 1 \rangle 4. \ S \neq \emptyset$ 

 $\langle 2 \rangle 1$ . PICK *n* such that  $a_n = a$ 

 $\langle 2 \rangle 2. \ b_n \in S$ 

 $\langle 1 \rangle 5$ . Let:  $s := \sup S$ 

 $\langle 1 \rangle 6. \ s \in S$ 

 $\langle 2 \rangle 1$ . PICK an increasing sequence  $(s_n)$  in S that converges to S.

 $\langle 2 \rangle 2$ . For all n we have  $s_n - a \leq \sum_{a_m < b_{sm}} (b_{sm} - a_m) \leq s - a$ . Proof:

$$s_n - a = \sum_{a_m < b_{s_n m}} (b_{s_n m} - a_m) \qquad (s_n \in S)$$

$$\leq \sum_{a_m < b_{sm}} (b_{sm} - a_m) \qquad (s_n \leq s)$$

$$\leq s - a \tag{\langle 1 \rangle 1}$$

 $\langle 2 \rangle 3.$   $\sum_{a_m < b_{sm}} (b_{sm} - a_m) = s - a$  Proof: Sandwich Theorem

 $\langle 1 \rangle 7. \ s = b$ 

 $\langle 2 \rangle 1$ . Assume: for a contradiction s < b

 $\langle 2 \rangle 2$ . Pick k such that  $s \in [a_k, b_k)$ 

 $\langle 2 \rangle 3. \ b_k \in S$ 

PROOF: Since  $\sum_{a_m < b_{sm}} (b_{sm} - a_m) = \sum_{a_m < b_{b,m}} (b_{b_k m} - a_m)$ .

 $\langle 2 \rangle 4. \ b_k \leq s$ 

 $\langle 2 \rangle 5$ . Q.E.D.

PROOF: This is a contradiction.

 $\langle 1 \rangle 8. \ b \in S$ 

$$\langle 1 \rangle 9. \sum_{a_n < b_n} (b_n - a_n) = b - a$$

**Theorem 17.1.17.** Let  $(f_n)$  be a decreasing sequence of non-negative step functions such that, for all  $x \in \mathbb{R}$ , we have  $f_n(x) \to 0$  as  $n \to \infty$ . Then  $\int f_n \to 0$ as  $n \to \infty$ .

#### Proof:

- $\langle 1 \rangle 1$ .  $(\int f_n)$  is decreasing and bounded below by 0.
- $\langle 1 \rangle 2$ . Let:  $\epsilon := \lim_{n \to 0} \int f_n$
- $\langle 1 \rangle 3$ . Assume: for a contradiction  $\epsilon > 0$ .
- $\langle 1 \rangle 4$ . Pick  $a, b \in \mathbb{R}$  such that supp  $f_0 \subseteq [a, b)$
- $\langle 1 \rangle 5$ . Let:  $\alpha := \epsilon/2(b-a)$
- $\langle 1 \rangle 6$ . For  $n \in \mathbb{N}$ , Let:

$$A_n := \{x \in [a,b) : f_n(x) < \alpha\}$$
.

- $\langle 1 \rangle 7$ . Let:  $B_0 := A_0$
- $\langle 1 \rangle 8$ . For *n* a positive integer, Let:  $B_n = A_n - A_{n-1}$ .
- $\langle 1 \rangle 9$ . For all n we have  $A_n \subseteq A_{n+1}$ .
- $\langle 1 \rangle 10$ . For  $m \neq n$  we have  $B_m \cap B_n = \emptyset$ .
- $\langle 1 \rangle 11. \bigcup_{n=0}^{\infty} A_n = [a, b)$

PROOF: For all  $x \in [a, b)$ , there exists N such that  $f_N(x) < \alpha$  because  $f_n(x) \to \infty$  $0 \text{ as } n \to \infty.$ 

- $\langle 1 \rangle 12. \ \bigcup_{n=0}^{\infty} B_n = [a, b)$
- $\langle 1 \rangle 13$ . For  $n \in \mathbb{N}$ ,

LET:  $B_n = [a_{n1}, b_{n1}) \cup \cdots \cup [a_{nk_n}, b_{nk_n}).$   $\langle 1 \rangle 14. \sum_{n=0}^{\infty} \sum_{k=1}^{k_n} (b_{nk} - a_{nk}) = b - a$ Proof: Lemma 17.1.16.

 $\langle 1 \rangle 15$ . Let:

$$\delta = \frac{\epsilon}{2 \max_{x} |f_0(x)|} .$$

 $\langle 1 \rangle 16$ . Pick  $n_0$  such that

$$\sum_{n=n_0}^{\infty} \sum_{k=1}^{k_n} (b_{nk} - a_{nk}) < \delta .$$

- $\langle 1 \rangle 17. \ A_{n_0} = B_0 \cup \cdots \cup B_{n_0}$
- $\langle 1 \rangle 18$ . Let:  $g: \mathbb{R} \to \mathbb{R}$ ,

$$g(x) = \begin{cases} f_{n_0}(x) & \text{if } x \in A_{n_0} \\ 0 & \text{if } x \notin A_{n_0} \end{cases}$$

 $\langle 1 \rangle 19$ . Let:  $h : \mathbb{R} \to \mathbb{R}$ ,

$$h(x) = \begin{cases} 0 & \text{if } x \in A_{n_0} \\ f_{n_0}(x) & \text{if } x \notin A_{n_0} \end{cases}$$

- $\langle 1 \rangle 20. \ \forall x \in B. f_{n_0}(x) < \alpha$
- $\langle 1 \rangle 21. \ \forall x \in \mathbb{R}. g(x) < \alpha$
- $\langle 1 \rangle 22$ .

$$\int g < \frac{\epsilon}{2}$$

 $\langle 1 \rangle 23.$ 

$$\int h < \frac{\epsilon}{2}$$

Proof:

$$\int h < \delta \max_{x} |f_{n_0}(x)|$$

$$\leq \delta \max_{x} |f_0(x)|$$

$$= \epsilon/2$$

 $\langle 1 \rangle 24$ .

$$\int f_{n_0} < \epsilon$$

 $\int f_{n_0} < \epsilon$  Proof: Since  $\forall x. f_{n_0}(x) = g(x) + h(x).$   $\rangle 25.$ 

 $\langle 1 \rangle 25$ .

$$\lim_{n \to \infty} \int f_n < \epsilon$$

 $\langle 1 \rangle 26$ . Q.E.D.

PROOF: This contradicts  $\langle 1 \rangle 2$ .

Corollary 17.1.17.1. Let  $(f_n)$  be an increasing sequence of step functions. If  $\lim_{n\to\infty} f_n(x) \ge 0$  for all x, then  $\lim_{n\to\infty} \int f_n \ge 0$ .

Proof:

 $\langle 1 \rangle 1$ . For  $n \in \mathbb{N}$ ,

Let:  $g_n : \mathbb{R} \to \mathbb{R}$  be the function  $g_n(x) = \max(0, -f_n(x))$ .

 $\langle 1 \rangle 2$ .  $(g_n)$  is a decreasing sequence of step functions and  $g_n(x) \to 0$  as  $n \to \infty$ for all x.

 $\langle 1 \rangle 3$ .  $\int g_n \to 0$  as  $n \to \infty$ 

 $\langle 1 \rangle 4$ . For all  $n \in \mathbb{N}$  we have  $f_n(x) = \max(0, f_n(x)) - \max(0, -f_n(x))$ .

 $\langle 1 \rangle 5$ .

$$\int f_n = \int \max(0, f_n(x)) dx - \int \max(0, -f_n(x)) dx$$

 $\langle 1 \rangle 6$ .

$$\langle 1 \rangle 6.$$
 
$$\lim_{n \to \infty} \int f_n = \lim_{n \to \infty} \int \max(0, f_n(x)) dx$$
 
$$\langle 1 \rangle 7. \lim_{n \to \infty} \int f_n \geq 0$$

#### 17.2Lebesgue Integration

**Definition 17.2.1** (Lebesgue Integration). Let  $f: \mathbb{R} \to \mathbb{R}$ . Then f is Lebesgue integrable iff there exists a sequence  $(f_n)$  of step functions such that:

$$\sum_{n=0}^{\infty} \int |f_n(x)| dx < \infty$$

• For all  $x \in \mathbb{R}$ , if  $\sum_{n=0}^{\infty} |f_n(x)| < \infty$ , then  $f(x) = \sum_{n=0}^{\infty} f_n(x)$ .

We write  $f \simeq \sum_{n=0}^{\infty} f_n$  for these two conditions. The *integral* of f is then

$$\int f := \sum_{n=0}^{\infty} \int f_n .$$

## Chapter 18

# Complex Analysis

**Theorem 18.0.1** (Hölder's Inequality). Let p and q be real numbers with p > 1, q > 1 and 1/p + 1/q = 1. If  $(x_n) \in l^p$  and  $(y_n) \in l^q$  then

$$\sum_{n=0}^{\infty} |x_n y_n| \le \left(\sum_{n=0}^{\infty} |x_n|^p\right)^{1/p} \left(\sum_{n=0}^{\infty} |y_n|^q\right)^{1/q}$$

Proof:

 $\langle 1 \rangle 1$ . Let: p and q be real numbers with p > 1 and q > 1

 $\langle 1 \rangle 2$ . Assume: 1/p + 1/q = 1

 $\langle 1 \rangle 3$ . Let:  $(x_n) \in l^p$ 

 $\langle 1 \rangle 4$ . Let:  $(y_n) \in l^q$ 

 $\langle 1 \rangle$ 5. Assume: w.l.o.g.  $x_0 \neq 0$  and  $y_0 \neq 0$ 

 $\langle 1 \rangle 6$ . For all  $x \in [0,1]$ , we have

$$x^{1/p} \le \frac{1}{p}x + \frac{1}{q} .$$

Proof:

 $\langle 2 \rangle 1$ . Let:  $f:[0,1] \to \mathbb{R}$  be the function

$$f(x) = \frac{1}{p}x + \frac{1}{q} - x^{1/p}$$
.

 $\langle 2 \rangle 2.$   $f'(x) = \frac{1}{p} - \frac{1}{p} x^{-1/q}$  for  $x \in (0, 1]$   $\langle 2 \rangle 3.$  f'(x) < 0 for  $x \in (0, 1]$ 

 $\langle 2 \rangle 4$ . f(1) = 1/p + 1/q - 1 = 0

 $\langle 2 \rangle 5$ .  $f(x) \geq 0$  for all  $x \in [0,1]$ 

 $\langle 1 \rangle$ 7. For all non-negative reals a and b, we have

$$ab \le \frac{a^p}{p} + \frac{b^q}{q} .$$

 $\langle 2 \rangle 1$ . Let: a and b be non-negative reals.

 $\langle 2 \rangle 2$ . Case:  $a^p \leq b^q$ 

 $\langle 3 \rangle 1. \ 0 \leq a^p/b^q \leq 1$ 

 $\langle 3 \rangle 2$ .

$$ab^{-q/p} \le \frac{1}{p} \frac{a^p}{b^q} + \frac{1}{q}$$

PROOF: Taking  $x = a^p/b^q$  in  $\langle 1 \rangle 6$ .

 $\langle 3 \rangle 3$ .

$$ab^{1-q} \le \frac{1}{p} \frac{a^p}{b^q} + \frac{1}{q}$$

PROOF: -q/p = 1 - q from  $\langle 1 \rangle 2$ 

 $\langle 3 \rangle 4$ .

$$ab \le \frac{a^p}{p} + \frac{b^q}{q}$$

 $\langle 2 \rangle 3$ . Case:  $b^q \leq a^p$ 

PROOF: Similar.

$$\langle 1 \rangle 8. \text{ For } j = 1, \dots, n, \text{ we have}$$

$$\frac{|x_j|}{\left(\sum_{k=0}^n |x_k|^p\right)^{1/p}} \frac{|y_j|}{\left(\sum_{k=0}^n |y_k|^q\right)^{1/q}} \le \frac{1}{p} \frac{|x_j|^p}{\sum_{k=0}^n |x_k|^p} + \frac{1}{q} \frac{|y_j|^q}{\sum_{k=0}^n |y_k|^q}$$

$$a = \frac{|x_j|}{\left(\sum_{k=0}^n |x_k|^p\right)^{1/p}}$$
 and  $b = \frac{|y_j|}{\left(\sum_{k=0}^n |y_k|^q\right)^{1/q}}$ .

 $\langle 1 \rangle 9$ .

$$\frac{\sum_{j=0}^{n} |x_j| |y_j|}{\left(\sum_{k=0}^{n} |x_k|^p\right)^{1/p} \left(\sum_{k=0}^{n} |y_k|^q\right)^{1/q}} \le 1$$

Proof:

FROOF: 
$$\frac{\sum_{j=0}^{n} |x_j| |y_j|}{\left(\sum_{k=0}^{n} |x_k|^p\right)^{1/p} \left(\sum_{k=0}^{n} |y_k|^q\right)^{1/q}} \le \frac{1}{p} + \frac{1}{q} \quad \text{(Taking the sum } j = 0 \text{ to } n \text{ in } \langle 1 \rangle 8)$$

$$= 1 \qquad \qquad (\langle 1 \rangle 2)$$

 $\langle 1 \rangle 10$ . Q.E.D.

PROOF: Taking the limit  $n \to \infty$  in  $\langle 1 \rangle 9$ .

**Theorem 18.0.2** (Minkowski's Inequality). Let p be a real number,  $p \ge 1$ . Let  $(x_n),(y_n)\in l^p$ . Then

$$\left(\sum_{n=1}^{\infty} |x_n + y_n|^p\right)^{1/p} \le \left(\sum_{n=1}^{\infty} |x_n|^p\right)^{1/p} + \left(\sum_{n=1}^{\infty} |y_n|^p\right)^{1/p}$$

Proof:

- $\langle 1 \rangle 1$ . Let: p be a real number with  $p \geq 1$
- $\langle 1 \rangle 2$ . Assume: w.l.o.g. p > 1

PROOF: The case p = 1 is just the Triangle Inequality.

- $\langle 1 \rangle 3$ . Let: q be the real such that 1/p + 1/q = 1
- $\langle 1 \rangle 4$ .

$$\sum_{n=0}^{\infty} |x_n + y_n|^p \le \left(\sum_{n=0}^{\infty} |x_n|^p\right)^{1/p} \left(\sum_{n=0}^{\infty} |x_n + y_n|^{q(p-1)}\right)^{1/q} + \left(\sum_{n=0}^{\infty} |y_n|^p\right)^{1/p} \left(\sum_{n=0}^{\infty} |x_n + y_n|^{q(p-1)}\right)^{1/q}$$

PROOF: 
$$\sum_{n=0}^{\infty} |x_n + y_n|^p = \sum_{n=0}^{\infty} |x_n + y_n| |x_n + y_n|^{p-1}$$

$$\leq \sum_{n=0}^{\infty} |x_n| |x_n + y_n|^{p-1} + \sum_{n=0}^{\infty} |y_n| |x_n + y_n|^{p-1}$$
 (Triangle Inequality) 
$$\leq \left( \sum_{n=0}^{\infty} |x_n|^p \right)^{1/p} \left( \sum_{n=0}^{\infty} |x_n + y_n|^{q(p-1)} \right)^{1/q}$$

$$+ \left( \sum_{n=0}^{\infty} |y_n|^p \right)^{1/p} \left( \sum_{n=0}^{\infty} |x_n + y_n|^{q(p-1)} \right)^{1/q}$$
 (Hölder's Inequality) 
$$\langle 1 \rangle 5.$$

$$\sum_{n=0}^{\infty} |x_n + y_n|^p \leq \left\{ \left( \sum_{n=0}^{\infty} |x_n|^p \right)^{1/p} + \left( \sum_{n=0}^{\infty} |y_n|^p \right)^{1/p} \right\} \left( \sum_{n=0}^{\infty} |x_n + y_n|^p \right)^{1/q}$$

 $\langle 1 \rangle 6$ . Q.E.D.

## Chapter 19

# Topology

## 19.1 Topological Spaces

**Definition 19.1.1** (Topology). Let X be a set. A topology on X is a set  $\mathcal{T} \subseteq \mathcal{P}X$ , whose elements are called *open sets*, such that:

- $X \in \mathcal{T}$
- $\forall \mathcal{U} \subseteq \mathcal{T}. \bigcup \mathcal{U} \in \mathcal{T}$
- $\forall U, V \in \mathcal{T}.U \cap V \in \mathcal{T}$

A topological space is a pair  $(X, \mathcal{T})$  such that X is a set and  $\mathcal{T}$  is a topology on X. We refer to the elements of X as points.

An open neighbourhood of a point x is an open set U such that  $x \in U$ . We write  $\mathcal{T}_x$  for the set of all open neighbourhoods of x.

**Definition 19.1.2** (Closed Set). In a topological space X, a set C is *closed* iff X - C is open.

**Definition 19.1.3** (Discrete Topology). The *discrete topology* on a set X is  $\mathcal{P}X$ .

**Definition 19.1.4** (Indiscrete Topology). The *indiscrete topology* or *trivial topology* on a set X is  $\{\emptyset, X\}$ .

**Definition 19.1.5** (Finer, Coarser). Let  $\mathcal{T}$  and  $\mathcal{T}'$  be topologies on the same set X. Then  $\mathcal{T}$  is *finer*, *larger* or *stronger* than  $\mathcal{T}'$ , and  $\mathcal{T}'$  is *coarser*, *smaller* or *weaker* than  $\mathcal{T}$ , iff  $\mathcal{T}' \subseteq \mathcal{T}$ .

**Definition 19.1.6** (Basis). Let X be a set. A *basis* for a topology on X is a set  $\mathcal{B} \subseteq \mathcal{P}X$ , whose elements we call *basic open neighbourhoods*, such that:

- $\bigcup \mathcal{B} = X$
- $\forall A, B \in \mathcal{B}. \forall x \in A \cap B. \exists C \in \mathcal{B}. x \in C \subseteq A \cap B.$

The topology generated by  $\mathcal{B}$  is the coarsest topology that includes  $\mathcal{B}$ .

**Proposition 19.1.7.** The topology generated by  $\mathcal{B}$  is  $\{U \in \mathcal{P}X \mid \forall x \in U.\exists B \in \mathcal{B}. x \in B \subseteq U\}$ 

### 19.2 Continuous Functions

**Definition 19.2.1** (Continuous). Let X and Y be topological spaces. Let  $f: X \to Y$ . Then f is *continuous* iff, for every open set U in Y, the set  $f^{-1}(U)$  is open in X.

**Proposition 19.2.2.** For any topological space X, the identity function  $id_X : X \to X$  is continuous.

**Proposition 19.2.3.** If  $f: X \to Y$  and  $g: Y \to Z$  are continuous then  $g \circ f: X \to Z$  is continuous.

**Definition 19.2.4.** Let Top be the category of topological spaces and continuous functions.

**Proposition 19.2.5.** Let X and Y be topological spaces and  $f: X \to Y$ . Then f is continuous if and only if, for every closed set C in Y, we have  $f^{-1}(C)$  is closed in X.

**Definition 19.2.6** (Continuous at a Point). Let X and Y be topological spaces. Let  $f: X \to Y$ . Let  $x \in X$ . Then f is *continuous at* x iff, for every open neighbourhood V of f(x), we have  $f^{-1}(V)$  is open.

**Definition 19.2.7.** The category of *pointed* topological spaces,  $Top_*$ , is the category with:

- objects all pairs (A, a) where A is a topological space and  $a \in A$ ;
- morphisms  $f:(A,a)\to (B,b)$  all continuous functions  $f:A\to B$  such that f(a)=b.

**Definition 19.2.8** (Homeomorphism). A homeomorphism is an isomorphism in Top. Two isomorphic topological spaces are called homeomorphic.

**Definition 19.2.9** (Topological Property). A property of topological spaces is a *topological* property iff it is preserved by homeomorphism.

**Proposition 19.2.10.** Cardinality of a topological space is a topological property.

## 19.3 Convergence

**Definition 19.3.1** (Convergence). Let X be a topological space. Let  $(x_n)$  be a sequence in X and  $l \in X$ . Then  $(x_n)$  converges to l,  $x_n \to l$  as  $n \to \infty$ , if and only if, for every open neighbourhood U of l, there exists N such that  $\forall n \geq N. x_n \in U$ .

**Theorem 19.3.2.** Let X and Y be topological spaces. Let Z be a closed subspace of X and  $f: Z \to Y$  a continuous function. Then the graph of f,  $G = \{(x, f(x)) \mid x \in Z\}$ , is closed in  $X \times Y$ .

Proof:

 $\langle 1 \rangle 1$ . Let:  $((x_n, f(x_n)))$  be a sequence in G.

 $\langle 1 \rangle 2$ . Let:  $(x_n, T(x_n)) \to (x, y)$  as  $n \to \infty$ 

 $\langle 1 \rangle 3. \ x \in W$ 

PROOF: Since  $x_n \to x$  and W is closed.

 $\langle 1 \rangle 4. \ y = f(x)$ 

Proof:

$$y = \lim_{n \to \infty} f(x_n)$$
$$= f\left(\lim_{n \to \infty} x_n\right)$$
$$= f(x)$$

## 19.4 Homotopy

**Definition 19.4.1.** Let hTop be the category whose objects are topological spaces, and whose morphisms are homotopy classes of continuous functions.

**Definition 19.4.2.** A *homotopy equivalence* is an isomorphism in hTop. Isomorphic topological spaces are called *homotopic*.

## 19.5 Metric Spaces

**Definition 19.5.1** (Metric). Let X be a set. A *metric* on a set X is a function  $d: X^2 \to \mathbb{R}$  such that:

- $\forall x, y \in X.d(x, y) \ge 0$
- $\forall x, y \in X.d(x, y) = d(y, x)$
- Triangle Inequality  $\forall x, y, z \in X.d(x, y) + d(y, z) \ge d(x, z)$
- $\forall x, y \in X.d(x, y) = 0 \text{ iff } x = y.$

A metric space is a pair (X, d) such that d is a metric on X.

**Definition 19.5.2** (Open Ball). In a metric space X, let  $c \in X$  and r > 0. The *open ball* with *centre* c and *radius* r is

$$B(c,r) := \{ x \in X \mid d(x,c) < r \}$$
.

**Proposition 19.5.3.** In a metric space, the set of open balls forms a basis for a topology.

**Definition 19.5.4** (Metric Topology). Given a metric space X, the *metric topology* on X is the topology generated by the basis of open balls.

A topological space  $(X, \mathcal{T})$  is *metrizable* iff there exists a metric d on X such that  $\mathcal{T}$  is the metric topology induced by d.

We identify a metric space with this topological space.

**Proposition 19.5.5.** If d is a metric on X and  $Y \subseteq X$  then  $d \upharpoonright Y^2$  is a metric on Y.

We write just Y for the metric space  $(Y, d \upharpoonright Y^2)$ .

**Proposition 19.5.6.** Let X and Y be metric spaces. Let  $f: X \to Y$ . Then f is continuous if and only if, for every sequence  $(x_n)$  in X and  $l \in X$ , if  $x_n \to l$  as  $n \to \infty$  then  $f(x_n) \to f(l)$  as  $n \to \infty$ .

#### Proof:

- $\langle 1 \rangle 1$ . If f is continuous then, for every sequence  $(x_n)$  in X and  $l \in X$ , if  $x_n \to l$  as  $n \to \infty$ , then  $f(x_n) \to f(l)$  as  $n \to \infty$ .
  - $\langle 2 \rangle 1$ . Assume: f is continuous.
  - $\langle 2 \rangle 2$ . Let:  $(x_n)$  be a sequence in X.
  - $\langle 2 \rangle 3$ . Let:  $l \in X$
  - $\langle 2 \rangle 4$ . Assume:  $x_n \to l$  as  $n \to \infty$ .
  - $\langle 2 \rangle$ 5. Let: V be an open neighbourhood of f(l)
  - $\langle 2 \rangle 6$ .  $f^{-1}(V)$  is an open neighbourhood of l.
  - $\langle 2 \rangle 7$ . PICK N such that  $\forall n \geq N.x_n \in f^{-1}(V)$
  - $\langle 2 \rangle 8. \ \forall n \geq N. f(x_n) \in V$
- $\langle 1 \rangle 2$ . If, for every sequence  $(x_n)$  in X and  $l \in X$ , if  $x_n \to l$  as  $n \to \infty$ , then  $f(x_n) \to f(l)$  as  $n \to \infty$ , then f is continuous.
  - $\langle 2 \rangle 1$ . Assume: f is not continuous.
  - $\langle 2 \rangle 2$ . PICK an open set V in Y such that  $f^{-1}(V)$  is not open in X
  - $\langle 2 \rangle 3$ . PICK  $l \in f^{-1}(V)$  such that, for all  $\epsilon > 0$ ,  $B(l, \epsilon) \nsubseteq f^{-1}(V)$ .
  - $\langle 2 \rangle 4$ . For  $n \in \mathbb{N}$ , Pick  $x_n \in B(l, 1/(n+1))$  such that  $x_n \notin f^{-1}(V)$ .
  - $\langle 2 \rangle 5$ .  $x_n \to l$  as  $n \to \infty$ .
  - $\langle 2 \rangle 6. \ f(x_n) \not\to f(l) \text{ as } n \to \infty$

**Proposition 19.5.7.** Completeness is not a topological property.

PROOF: We have  $(-1,1) \cong \mathbb{R}$ , but  $\mathbb{R}$  is complete and (-1,1) is not.  $\sqcup$ 

# Chapter 20

# Ring Theory

**Definition 20.0.1.** Given a ring R, let R – Mod be the category of modules over R and R-linear maps.

## Chapter 21

# Linear Algebra

## 21.1 Vector Spaces

**Definition 21.1.1** (Vector Space). Let K be a field. A *vector space* over K consists of:

- a set V, whose elements are called *vectors*;
- an operation  $+: V^2 \to V$ , addition;
- an operation  $\cdot: K \times V \to V$ , scalar multiplication

such that:

- V is an Abelian group under +
- $\forall \alpha, \beta \in K. \forall x \in V. \alpha(\beta x) = (\alpha \beta) x$
- $\forall \alpha, \beta \in K. \forall x \in V. (\alpha + \beta)x = \alpha x + \beta x$
- $\forall \alpha \in K. \forall x, y \in V. \alpha(x+y) = \alpha x + \alpha y$
- $\forall x \in V.1x = x$

We call the elements of K scalars. A real vector space is a vector space over  $\mathbb{R}$ , and a complex vector space is a vector space over  $\mathbb{C}$ .

**Proposition 21.1.2.** *Let* K *be a field. Let* V *be a vector space over* K. *For any*  $\lambda \in K$  *we have*  $\lambda 0 = 0$ .

$$\lambda 0 = \lambda(0+0)$$

$$= \lambda 0 + \lambda 0$$

$$\therefore 0 = \lambda 0$$

**Proposition 21.1.3.** Let K be a field. Let V be a vector space over K. Let  $\lambda \in K$  and  $x \in V$ . If  $\lambda x = 0$  then either  $\lambda = 0$  or x = 0.

PROOF: If  $\lambda \neq 0$  then  $x = 1x = \lambda^{-1}\lambda x = \lambda^{-1}0 = 0$ .

**Proposition 21.1.4.** Let K be a field. Let V be a vector space over K. For any  $x \in V$  we have 0x = 0.

Proof:

$$0x = (0+0)x$$
$$= 0x + 0x$$
$$\therefore 0 = 0x$$

**Proposition 21.1.5.** Let K be a field. Let V be a vector space over K. For any  $x \in V$ , we have (-1)x = -x.

Proof:

$$x + (-1)x = 1x + (-1)x$$

$$= (1 + (-1))x$$

$$= 0x$$

$$= 0$$

$$\therefore (-1)x = -x$$

**Proposition 21.1.6.** Let K be a field. Then K is a vector space over K under addition and multiplication in K.

Proof: Easy.  $\square$ 

**Proposition 21.1.7.**  $\mathbb{C}$  *is a vector space over*  $\mathbb{R}$ .

Proof: Easy.  $\square$ 

**Proposition 21.1.8.** Let K be a field. Let  $\{V_i\}_{i\in I}$  be a family of vector spaces over K. Then  $\prod_{i\in I} V_i$  is a vector space under

$$(f+g)(i) = f(i) + g(i) \qquad (f,g \in \prod_{i \in I} V_i, x \in X)$$
$$(\lambda f)(x) = \lambda f(x) \qquad (\lambda \in K, f \in \prod_{i \in I} V_i, x \in X)$$

Proof: Easy.  $\square$ 

## 21.2 Subspaces

**Definition 21.2.1** (Vector Subspace). Let K be a field. Let V be a vector space over K. A *vector subspace* of V is a subset  $U \subseteq V$  such that, for all  $\alpha, \beta \in K$  and  $x, y \in U$ , we have  $\alpha x + \beta y \in U$ .

It is a proper subspace iff  $U \neq V$ .

**Proposition 21.2.2.** If U is a subspace of V then U is a vector space under the restrictions of + and  $\cdot$  to U.

PROOF: Easy.

**Proposition 21.2.3.** V is a subspace of V.

Proof: Easy.

**Proposition 21.2.4.** If U is a subspace of V and V is a subspace of W then U is a subspace of W.

Proof: Easy.

**Definition 21.2.5.** Let  $\Omega$  be a topological space. Then  $\mathcal{C}(\Omega)$  is the complex vector space of all continuous functions from  $\Omega$  to  $\mathbb{C}$ . This is a subspace of  $\mathbb{C}^{\Omega}$ .

**Definition 21.2.6.** Let  $n, k \in \mathbb{N}$ . Let  $\Omega$  be an open subset of  $\mathbb{R}^n$ . Then  $\mathcal{C}^k(\Omega)$  is the complex vector space of all functions  $\Omega \to \mathbb{C}$  that have all continuous partial derivatives of order k. This is a subspace of  $\mathcal{C}(\Omega)$ . If l > k then  $\mathcal{C}^l(\Omega)$  is a subpase of  $\mathcal{C}^k(\Omega)$ .

**Definition 21.2.7.** Let  $n \in \mathbb{N}$ . Let  $\Omega$  be an open subset of  $\mathbb{R}^n$ . Then  $\mathcal{C}^{\infty}(\Omega)$  is the complex vector space of all infinitely differentiable functions  $\Omega \to \mathbb{C}$ . This is a subspace of  $\mathcal{C}^k(\Omega)$  for all k.

**Definition 21.2.8.** Let  $n \in \mathbb{N}$ . Let  $\Omega$  be an open subset of  $\mathbb{R}^n$ . Then  $\mathcal{P}(\Omega)$  is the complex vector space of all complex polynomials of n variables, considered as functions  $\Omega \to \mathbb{C}$ . This is a subspace of  $\mathcal{C}^{\infty}(\Omega)$ .

**Proposition 21.2.9.** The space of all convergent sequences in  $\mathbb{C}$  is a subspace of the space of all bounded sequences in  $\mathbb{C}$ , which is a subspace of  $\mathbb{C}^{\mathbb{N}}$ .

Proof: Easy.

**Definition 21.2.10.** Let p be a real number,  $p \ge 1$ . Let  $l^p$  be the set of all complex sequences  $(z_n)$  such that  $\sum_{n=1}^{\infty} |z_n|^p < \infty$ .

**Proposition 21.2.11.** For p a real number  $\geq 1$ , we have that  $l^p$  is a subspace of  $\mathbb{C}^{\mathbb{N}}$ .

#### Proof:

 $\langle 1 \rangle 1$ . For all  $(x_n), (y_n) \in l^p$ , we have  $(x_n + y_n) \in l^p$ .

PROOF: From Minkowski's Inequality.

 $\langle 1 \rangle 2$ . For all  $\lambda \in \mathbb{C}$  and  $(x_n) \in l^p$  we have  $(\lambda x_n) \in l^p$  PROOF:

$$\sum_{n=1}^{\infty} |\lambda x_n|^p = |\lambda|^p \sum_{n=1}^{\infty} |x_n|^p < \infty$$

**Definition 21.2.12** (Linear Combination). Let K be a field. Let V be a vector space over K. Let  $x, x_1, \ldots, x_n \in V$ . Then x is a linear combination of  $x_1, \ldots, x_n$  iff there exist  $\alpha_1, \ldots, \alpha_n \in K$  such that

$$x = \alpha_1 x_1 + \dots + \alpha_n x_n .$$

**Definition 21.2.13** (Linearly Independent). A finite set of vectors  $\{x_1, \ldots, x_n\}$  is *linearly independent* iff, whenever  $\alpha_1 x_1 + \cdots + \alpha_n x_n = 0$ , then  $\alpha_1 = \cdots = \alpha_n = 0$ .

A set of vectors is *linearly independent* iff every finite subset is linearly independent; otherwise, it is *linearly dependent*.

**Definition 21.2.14** (Span). Let  $\mathcal{A}$  be a set of vectors. The *span* of  $\mathcal{A}$ , span  $\mathcal{A}$ , is the set of all linear combinations of elements of  $\mathcal{A}$ .

**Proposition 21.2.15.** span A is the smallest subspace of V that includes A.

Proof: Easy.

**Definition 21.2.16** (Basis). A *basis* for V is a linearly independent set of vectors  $\mathcal{B}$  such that span  $\mathcal{B} = V$ .

**Definition 21.2.17** (Finite Dimensional). A vector space is *finite dimensional* iff it has a finite basis; otherwise it is *infinite dimensional*.

**Proposition 21.2.18.** In a finite dimensional vector space, any two bases have the same number of elements.

**Definition 21.2.19** (Dimension). The *dimension* of a finite dimensional vector space V, dim V, is the number of elements in any basis.

Proposition 21.2.20.

$$\dim K^n = n$$

PROOF: The standard basis is the set of vectors with one coordinate 1 and all others 0.  $\Box$ 

**Proposition 21.2.21.** The dimension of  $\mathbb{C}^n$  as a real vector space is 2n.

**Proposition 21.2.22.** The set of all step functions is a subspace of  $\mathbb{R}^{\mathbb{R}}$ .

### 21.3 Linear Transformations

**Definition 21.3.1** (Linear Transformation). Let K be a field. Let U and V be vector spaces over K. Let  $T: U \to V$ . Then T is a linear transformation iff

$$\forall \alpha, \beta \in K. \forall x, y \in U. T(\alpha x + \beta y) = \alpha T(x) + \beta T(y)$$
.

Let  $\mathrm{Vect}_K$  be the category of vector spaces over K and linear transformations.

**Proposition 21.3.2.** If  $T: U \to V$  is a linear transformation then T(U) is a subspace of V.

**Proposition 21.3.3.** If  $T: U \to V$  is a linear transformation then the graph of T,  $\{(x, T(x)) \mid x \in U\}$ , is a subspace of  $U \times V$ .

**Definition 21.3.4** (Null Space). Let U and V be vector spaces over K and  $T:U\to V$ . The *null space* of T is

$$\mathcal{N}(T) := \{ x \in U \mid T(x) = 0 \} .$$

**Proposition 21.3.5.** If  $T: U \to V$  is a linear transformation then  $\mathcal{N}(T)$  is a subspace of U.

**Proposition 21.3.6.** Let U and V be vector spaces over K. The set of all linear transformations  $U \to V$  is a vector space over K under

$$(S+T)(u) = S(u) + T(u)$$
$$(\lambda S)(u) = \lambda S(u)$$

### 21.4 Normed Spaces

**Definition 21.4.1** (Norm). Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . A *norm* on a vector space V over K is a function  $\| \ \| : V \to \mathbb{R}$  such that:

- 1.  $\forall x \in V ||x|| = 0 \Rightarrow x = 0$
- 2.  $\forall \lambda \in K. \forall x \in V. ||\lambda x|| = |\lambda| ||x||$
- 3. Triangle Inequality  $\forall x, y \in V ||x + y|| \le ||x|| + ||y||$

**Proposition 21.4.2.** Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . Let V be a normed space over K. Define  $d: V^2 \to \mathbb{R}$  by  $d(x,y) = \|x-y\|$ . Then d is a metric on V.

Proof: Easy.

We identify any normed space V with this metric space.

**Proposition 21.4.3.** *Let* K *be either*  $\mathbb{R}$  *or*  $\mathbb{C}$ . *Let* V *be a vector space over* K. *Let*  $\| \ \|$  *be a norm on* V. *Then* 

$$||0|| = 0$$
.

Proof:

$$||0|| = ||0 \cdot 0||$$
 (Proposition 21.1.4)  
=  $|0|||0||$  (Axiom 2 for a norm)  
= 0

**Proposition 21.4.4.** *Let* K *be either*  $\mathbb{R}$  *or*  $\mathbb{C}$ . *Let* V *be a vector space over* K. *Let*  $\parallel \parallel$  *be a norm on* V. *Let*  $x \in V$ . *Then* 

$$||x|| \ge 0$$
.

Proof:

$$0 = ||0||$$
 (Proposition 21.4.3)  

$$= ||x - x||$$
 (Triangle Inequality)  

$$= ||x|| + ||x||$$
 (Axiom 2 for a norm)  

$$= 2||x||$$

**Proposition 21.4.5.** Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . Let V be a vector space over K. Let  $\| \cdot \|$  be a norm on V. Let  $x, y \in V$ . Then

$$|||x|| - ||y||| \le ||x - y||$$
.

Proof:

 $\langle 1 \rangle 1. ||x|| - ||y|| \le ||x - y||$ 

PROOF:  $||x|| \le ||x - y|| + ||y||$  by the Triangle Inequality.

 $\langle 1 \rangle 2$ .  $||y|| - ||x|| \le ||x - y||$ 

Proof:

$$||x|| + ||x - y|| = ||x|| + ||y - x||$$
 (Axiom 2 for a norm)  
 $\leq ||y||$  (Triangle Inequality)

**Corollary 21.4.5.1.** *Let* V *be a normed space. Then*  $\| \| : V \to \mathbb{R}$  *is continuous.* 

**Definition 21.4.6** (Euclidean Norm). The *Euclidean norm* on  $\mathbb{C}^n$  is defined by

$$||(z_1,\ldots,z_n)|| = \sqrt{|z_1|^2 + \cdots + |z_n|^2}$$

**Proposition 21.4.7.** Define  $\| \| : \mathbb{C}^n \to \mathbb{R}$  by

$$||(z_1,\ldots,z_n)|| = |z_1| + \cdots + |z_n|$$

Then this defines a norm on  $\mathbb{C}^n$ .

Proof: Easy.  $\square$ 

**Proposition 21.4.8.** Define  $\| \cdot \| : \mathbb{C}^n \to \mathbb{R}$  by

$$||(z_1,\ldots,z_n)|| = \max(|z_1|,\ldots,|z_n|)$$

Then this defines a norm on  $\mathbb{C}^n$ .

Proof: Easy.  $\square$ 

**Proposition 21.4.9.** Let  $\Omega$  be a closed bounded subset of  $\mathbb{R}^n$ . Define  $\| \ \| : \mathcal{C}(\Omega) \to \mathbb{R}$  by  $\|f\| = \max_{x \in \Omega} |f(x)|$ . Then  $\| \ \|$  defines a norm on  $\mathcal{C}(\Omega)$ .

Proof: Easy.  $\square$ 

**Proposition 21.4.10.** Let p be a real number,  $p \ge 1$ . Define  $\| \| : l^p \to \mathbb{R}$  by

$$||(z_n)|| = \left(\sum_{n=0}^{\infty} |z_n|^p\right)^{1/p}$$
.

Then this defines a norm on  $l^p$ .

PROOF: Easy. The triangle inequality is Minkowski's Inequality.  $\square$ 

**Definition 21.4.11** (Normed Space). Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . A normed space over K consists of a vector space V over K and a norm on V.

We shall write simply:

- $K^n$  for the normed space  $K^n$  under the Euclidean norm
- $l^p$  for the normed space  $l^p$  under the norm  $||(z_n)|| = (\sum_{n=0}^{\infty} |z_n|^p)^{1/p}$ .

**Proposition 21.4.12.** Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . Let V be a normed space over K. If  $x_n \to l$  as  $n \to \infty$  in V and  $\lambda_n \to \lambda$  as  $n \to \infty$  in K, then  $\lambda_n x_n \to \lambda l$  as  $n \to \infty$ .

Proof:

 $\langle 1 \rangle 1$ . Let:  $\epsilon > 0$ 

 $\langle 1 \rangle 2$ . Let:  $K = |\lambda| + \epsilon/2||l||$ 

(1)3. PICK N such that, for all  $n \ge N$ , we have  $|\lambda_n - \lambda| < \epsilon/2||l||$  and  $||x_n - l|| < \epsilon/(2K)$ 

 $\langle 1 \rangle 4$ . For all  $n \geq N$  we have  $|\lambda_n| < K$ 

 $\langle 1 \rangle 5. \|\lambda_n x_n - \lambda l\| < \epsilon$ 

Proof:

$$\begin{aligned} \|\lambda_n x_n - \lambda l\| &\leq \|\lambda_n x_n - \lambda_n l\| + \|\lambda_n l - \lambda l\| \\ &= |\lambda_n| \|x_n - l\| + |\lambda_n - \lambda| \|l\| \\ &< K \frac{\epsilon}{2K} + \frac{\epsilon}{2\|l\|} \|l\| \\ &= \epsilon \end{aligned}$$

**Proposition 21.4.13.** In a normed space, if  $x_n \to l$  and  $y_n \to m$  then  $x_n + y_n \to l + m$ 

Proof:

$$||(x_n + y_n) - (l+m)|| \le ||x_n - l|| + ||y_n - m||$$
  
 $\to 0$ 

**Definition 21.4.14** (Uniform Convergence). Let  $\Omega$  be a closed bounded set in  $\mathbb{R}^n$ . Let  $(f_n)$  be a sequence in  $\mathcal{C}(\Omega)$  and  $f \in \mathcal{C}(\Omega)$ . Then  $(f_n)$  converges uniformly to f if and only if, for every  $\epsilon > 0$ , there exists N such that  $\forall x \in \Omega. \forall n \geq N. |f_n(x) - f(x)| < \epsilon$ .

**Proposition 21.4.15.**  $(f_n)$  converges uniformly to f iff  $(f_n)$  converges to f under the uniform convergence norm.

Proof: Easy.

**Proposition 21.4.16.** There is no norm on C([0,1]) that induces pointwise convergence.

### Proof:

 $\langle 1 \rangle 1$ . Let:  $\| \|$  be any norm on  $\mathcal{C}([0,1])$ 

 $\langle 1 \rangle 2$ . For  $n \in \mathbb{N}$ , define  $g_n \in \mathcal{C}([0,1])$  by

$$g_n(t) = \begin{cases} 2^n t & \text{if } 0 \le t \le 2^{-n} \\ 2 - 2^n t & \text{if } 2^{-n} \le t \le 2^{1-n} \\ 0 & \text{otherwise} \end{cases}$$

- $\langle 1 \rangle 3$ . For all  $n \in \mathbb{N}$  we have  $||g_n|| \neq 0$
- $\langle 1 \rangle 4$ . For  $n \in \mathbb{N}$ ,

LET:  $f_n = g_n / \|g_n\|$ 

- $\langle 1 \rangle 5$ . For all  $n \in \mathbb{N}$ ,  $||f_n|| = 1$
- $\langle 1 \rangle 6$ .  $f_n$  does not converge to 0
- $\langle 1 \rangle 7$ .  $f_n \to 0$  as  $n \to \infty$  pointwise.

**Definition 21.4.17** (Equivalent Norms). Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . Let V be a vector spaces over K. Then two norms  $\| \|_1$  and  $\| \|_2$  are *equivalent* if and only if, for any sequence  $(x_n)$  in V and  $l \in V$ , we have  $x_n \to l$  under  $\| \|_1$  if and only if  $x_n \to l$  under  $\| \|_2$ .

**Proposition 21.4.18.** Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . Let V be a vector spaces over K. Let  $\| \ \|_1$  and  $\| \ \|_2$  be norms on V. Then  $\| \ \|_1$  and  $\| \ \|_2$  are equivalent if and only if there exist positive reals  $\alpha$  and  $\beta$  such that, for all  $x \in V$ ,

$$\alpha \|x\|_1 \le \|x\|_2 \le \beta \|x\|_1 \tag{21.1}$$

#### Proof:

- $\langle 1 \rangle 1$ . If  $\| \|_1$  and  $\| \|_2$  are equivalent then (21.1) holds.
  - $\langle 2 \rangle 1$ . Assume:  $\| \|_1$  and  $\| \|_2$  are equivalent.
  - $\langle 2 \rangle 2$ . There exists  $\alpha > 0$  such that, for all  $x \in V$ , we have  $\alpha \|x\|_1 \leq \|x\|_2$ 
    - (3)1. Assume: for a contradiction  $\forall \alpha > 0. \exists x \in V.\alpha ||x||_1 > ||x||_2$
    - $\langle 3 \rangle 2$ . For  $n \in \mathbb{Z}^+$ , choose  $x_n \in V$  such that  $1/n \|x_n\|_1 > \|x_n\|_2$
    - $\langle 3 \rangle 3$ . For  $n \in \mathbb{Z}^+$ ,

Let:

$$y_n = \frac{1}{\sqrt{n}} \frac{x_n}{\|x_n\|_2}$$

- $\langle 3 \rangle 4$ .  $||y_n||_2 \to 0$  as  $n \to \infty$
- $\langle 3 \rangle 5$ . For all  $n \in \mathbb{Z}^+$ ,  $||y_n||_1 > \sqrt{n}$
- $\langle 3 \rangle 6$ .  $||y_n|| \not\to 0$  as  $n \to \infty$
- $\langle 2 \rangle$ 3. There exists  $\beta > 0$  such that, for all  $x \in V$ , we have  $||x||_2 \le \beta ||x||_1$  PROOF: Similar.

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\langle 1 \rangle 2. If (21.1) holds then \| \|_1 and \| \|_2 are equivalent.
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- $\langle 2 \rangle 1$ . Assume: (21.1) holds.
- $\langle 2 \rangle 2$ . Let:  $(x_n)$  be a sequence in V and  $l \in V$
- $\langle 2 \rangle 3$ . If  $x_n \to l$  under  $\| \|_1$  then  $x_n \to l$  under  $\| \|_2$ .
  - $\langle 3 \rangle 1$ . Assume:  $x_n \to l$  und  $\| \|_1$ .
  - $\langle 3 \rangle 2$ . Let:  $\epsilon > 0$
  - $\langle 3 \rangle 3$ . PICK N such that  $\forall n \geq N . ||x_n l|| < \epsilon/\beta$
  - $\langle 3 \rangle 4$ . Let:  $n \geq N$
  - $\langle 3 \rangle 5. \|x_n l\|_2 < \epsilon$

Proof:

$$||x_n - l||_2 \le \beta ||x_n - l||_1$$

 $<\epsilon$ 

 $\langle 2 \rangle 4$ . If  $x_n \to l$  under  $\| \|_2$  then  $x_n \to l$  under  $\| \|_1$ .

PROOF: Similar.

**Proposition 21.4.19.** Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . Let V be a normed space over K. If  $x_1, \ldots, x_n \in V$  are linearly independent, then there exists c > 0 such that, for all  $\alpha_1, \ldots, \alpha_n \in K$ ,

$$\|\alpha_1 x_1 + \dots + \alpha_n x_n\| \ge c(|\alpha_1| + \dots + |\alpha_n|)$$
.

Proof:

- $\langle 1 \rangle 1$ . Let:  $B = \{ (\beta_1, \dots, \beta_n) \in K^n \mid |\beta_1| + \dots + |\beta_n| = 1 \}$
- $\langle 1 \rangle 2$ . Let:  $f: B \to \mathbb{R}$  be the function

$$f(\beta_1,\ldots,\beta_n) = \|\beta_1 x_1 + \cdots + \beta_n x_n\|.$$

 $\langle 1 \rangle 3$ . Let: c be the minimum value in f(B)

PROOF: f is continuous and B is compact.

 $\langle 1 \rangle 4. \ c > 0$ 

PROOF: We never have  $f(\beta_1, \ldots, \beta_n) = 0$  by linear independence.

- $\langle 1 \rangle 5$ . Let:  $\alpha_1, \ldots, \alpha_n \in K$
- $\langle 1 \rangle$ 6. Assume: w.l.o.g.  $\alpha_1, \ldots, \alpha_n$  are not all zero.
- $\langle 1 \rangle 7$ . For  $i = 1, \ldots, n$ ,

Let: 
$$\beta_i = \alpha_i/(|\alpha_1| + \cdots + |\alpha_n|)$$

- $\langle 1 \rangle 8. \ (\beta_1, \ldots, \beta_n) \in B$
- $\langle 1 \rangle 9. \ f(\beta_1, \dots, \beta_n) \ge c$
- $\langle 1 \rangle 10. \|\alpha_1 x_1 + \dots + \alpha_n x_n\| \ge c(|\alpha_1| + \dots + |\alpha_n|)$

**Theorem 21.4.20.** Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . Let V be a finite dimensional vector space over K. Then any two norms on V are equivalent.

Proof:

- $\langle 1 \rangle 1$ . PICK a basis  $\{e_1, \ldots, e_n\}$  for V
- $\langle 1 \rangle 2$ . Let:  $\| \|_0 : V \to \mathbb{R}$  be the function

$$\|\alpha_1 e_1 + \dots + \alpha_n e_n\|_0 = |\alpha_1| + \dots + |\alpha_n|$$
.

 $\langle 1 \rangle 3$ .  $\| \|_0$  is a norm.

PROOF: If 
$$|\alpha_1| + \cdots + |\alpha_n| = 0$$
 then  $\alpha_1 = \cdots = \alpha_n = 0$ .  
 $\langle 2 \rangle 2$ .  $\forall \lambda \in K. \forall x \in V. ||\lambda x|| = |\lambda| ||x||$   
PROOF: 
$$||\lambda(\alpha_1 e_1 + \cdots + \alpha_n e_n)||_0 = ||\lambda \alpha_1 e_1 + \cdots + \lambda \alpha_n e_n||_0$$

$$= |\lambda \alpha_1| + \cdots + |\lambda \alpha_n|$$

$$= |\lambda|(|\alpha_1| + \cdots + |\alpha_n|)$$

$$= |\lambda|||\alpha_1 e_1 + \cdots + \alpha_n e_n||$$

 $\langle 2 \rangle 3$ . The triangle inequality holds.

 $\langle 2 \rangle 1. \ \forall x \in V. ||x||_0 = 0 \Rightarrow x = 0$ 

Proof:

$$\begin{aligned} \|(\alpha_1 e_1 + \dots + \alpha_n e_n) + (\beta_1 e_1 + \dots + \beta_n e_n)\| &= \|(\alpha_1 + \beta_1) e_1 + \dots + (\alpha_n + \beta_n) e_n\| \\ &= |\alpha_1 + \beta_1| + \dots + |\alpha_n + \beta_n| \\ &\leq (|\alpha_1| + \dots + |\alpha_n|) + (|\beta_1| + \dots + |\beta_n|) \\ &= \|\alpha_1 e_1 + \dots + \alpha_n e_n\|_0 + \|\beta_1 e_1 + \dots + \beta_n e_n\|_0 \end{aligned}$$

 $\langle 1 \rangle 4$ . Let:  $\| \|$  be any norm on V.

Prove:  $\| \|$  is equivalent to  $\| \|_0$ .

 $\langle 1 \rangle 5$ . For all  $\alpha_1, \ldots, \alpha_n \in K$ ,

$$\|\alpha_1 e_1 + \dots + \alpha_n e_n\| \le \max(\|e_1\|, \dots, \|e_n\|)(|\alpha_1| + \dots + |\alpha_n|)$$

 $\langle 2 \rangle 1$ . Let:  $\alpha_1, \ldots, \alpha_n \in K$ 

$$\langle 2 \rangle 2$$
.  $\|\alpha_1 e_1 + \dots + \alpha_n e_n\| \le \max(\|e_1\|, \dots, \|e_n\|)(|\alpha_1| + \dots + |\alpha_n|)$   
PROOF:

$$\|\alpha_1 e_1 + \dots + \alpha_n e_n\| \le |\alpha_1| \|e_1\| + \dots + |\alpha_n| \|e_n\|$$
  
  $\le (|\alpha_1| + \dots + |\alpha_n|) \max(\|e_1\|, \dots, \|e_n\|)$ 

- $\langle 1 \rangle 6$ . Let:  $\beta = \max(\|e_1\|, \dots, \|e_n\|)$
- $\langle 1 \rangle 7$ . For all  $x \in V$ ,

$$||x|| \le \beta ||x||_0 .$$

 $\langle 1 \rangle 8$ . There exists  $\alpha > 0$  such that, for all  $x \in V$ ,

$$\alpha \|x\|_0 \le \|x\| .$$

Proof: Proposition 21.4.19.

**Definition 21.4.21** (Closed Ball). Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . Let V be a vector space over K. Let  $x \in V$  and let r be a positive real number. The *closed ball* with *centre* x and *radius* r is

$$\overline{B(x,r)}:=\{y\in V\mid \|x-y\|\leq r\}\ .$$

**Definition 21.4.22** (Sphere). Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . Let V be a vector space over K. Let  $x \in V$  and let r be a positive real number. The *sphere* with *centre* x and radius r is

$$S(x,r) := \{ y \in V \mid ||x - y|| = r \} .$$

Proposition 21.4.23. Every closed ball is closed.

Proposition 21.4.24. Every sphere is closed.

Proposition 21.4.25. The union of two closed sets is closed.

**Proposition 21.4.26.** The intersection of a nonempty set of closed sets is closed.

**Proposition 21.4.27.** In a normed space V, both  $\emptyset$  and V are closed.

**Proposition 21.4.28.** Let V be a normed space and  $C \subseteq V$ . Then C is closed iff, for every sequence  $(x_n)$  in C and  $l \in V$ , if  $x_n \to l$  then  $l \in C$ .

# Proof:

- $\langle 1 \rangle 1$ . If C is closed then, for every sequence  $(x_n)$  in C and  $l \in V$ , if  $x_n \to l$  then  $l \in C$ .
  - $\langle 2 \rangle 1$ . Assume: C is closed.
  - $\langle 2 \rangle 2$ . Let:  $(x_n)$  be a sequence in C.
  - $\langle 2 \rangle 3$ . Let:  $l \in V$
  - $\langle 2 \rangle 4$ . Assume:  $x_n \to l$  as  $n \to \infty$
  - $\langle 2 \rangle$ 5. Assume: for a contradiction  $l \notin C$
  - $\langle 2 \rangle 6$ . PICK  $\epsilon > 0$  such that  $B(l, \epsilon) \subseteq V C$
  - $\langle 2 \rangle$ 7. PICK N such that  $\forall n \geq N . ||x_n l|| < \epsilon$
  - $\langle 2 \rangle 8. \ x_N \in C$
  - $\langle 2 \rangle 9. \|x_N l\| < \epsilon$
  - $\langle 2 \rangle 10$ . Q.E.D.

PROOF: This is a contradiction.

- $\langle 1 \rangle 2$ . If, for every sequence  $(x_n)$  in C and  $l \in V$ , if  $x_n \to l$  then  $l \in C$ , then C is closed.
  - $\langle 2 \rangle 1$ . Assume: for every sequence  $(x_n)$  in C and  $l \in V$ , if  $x_n \to l$  then  $l \in C$ .
  - $\langle 2 \rangle 2$ . Let:  $x \in V C$
  - $\langle 2 \rangle 3$ . Assume: for a contradiction there is no  $\epsilon > 0$  such that  $B(x,\epsilon) \subseteq V C$
  - $\langle 2 \rangle 4$ . For  $n \in \mathbb{Z}^+$ , Pick  $x_n \in B(x, \epsilon) \cap C$
  - $\langle 2 \rangle 5$ .  $x_n \to x$  as  $n \to \infty$
  - $\langle 2 \rangle 6$ . Q.E.D.

Proof: This is a contradiction.

**Definition 21.4.29** (Closure). Let V be a normed space and  $A \subseteq V$ . The *closure* of A, cl A, is the intersection of all the closed sets that include A.

**Theorem 21.4.30.** Let V be a normed space and  $S \subseteq V$ . Then the closure of S is the set of all limits of convergent sequences in S.

### Proof:

- $\langle 1 \rangle 1$ . For all  $l \in \operatorname{cl} S$ , there exists a sequence  $(x_n)$  that converges to l.
  - $\langle 2 \rangle 1$ . Let:  $l \in \operatorname{cl} S$
  - $\langle 2 \rangle 2$ . For  $n \in \mathbb{N}$ , Pick  $x_n \in B(l, 1/(n+1)) \cap S$ 
    - $\langle 3 \rangle 1$ . Assume: for a contradiction B(l, 1/(n+1)) does not intersect S.
    - $\langle 3 \rangle 2$ . V B(l, 1/(n+1)) is a closed set that includes S.
    - $\langle 3 \rangle 3$ . cl  $S \subseteq V B(l, 1/(n+1))$

```
\langle 3 \rangle 4. \ l \notin B(l, 1/(n+1))

\langle 3 \rangle 5. \ \text{Q.E.D.}

PROOF: This is a contradiction.

\langle 2 \rangle 3. \ x_n \to l

\langle 1 \rangle 2. For every sequence (x_n) in S, if x_n \to l then l \in \text{cl } S.

PROOF: Proposition 21.4.28.
```

**Definition 21.4.31** (Dense). Let V be a normed space and  $S \subseteq V$ . Then S is dense iff  $\operatorname{cl} S = V$ .

**Proposition 21.4.32.** In C([a,b]), the set of polynomials is dense.

PROOF: By the Weierstrass Theorem.  $\square$ 

**Proposition 21.4.33.** For any real  $p \ge 1$ , the set of all sequences with only finitely many non-zero terms is dense in  $l^p$ .

#### Proof:

 $\begin{array}{ll} \langle 1 \rangle 1. & \text{Let: } p \geq 1 \\ \langle 1 \rangle 2. & \text{Let: } (z_n) \in l^p \\ \langle 1 \rangle 3. & \text{Let: } \epsilon > 0 \\ \langle 1 \rangle 4. & \text{Pick } N \text{ such that } \forall n \geq N. |z_n| < \epsilon/2 \\ \langle 1 \rangle 5. & \text{Let: } (y_n) \text{ be the sequence with } y_n = z_n \text{ for } n < N, \text{ and } y_n = 0 \text{ for } n \geq N \\ \langle 1 \rangle 6. & \|(z_n) - (y_n)\| \leq \epsilon/2 \\ & \underline{\langle 1 \rangle} 7. & \|(z_n) - (y_n)\| < \epsilon \end{array}$ 

**Theorem 21.4.34.** Let V be a normed space. Let  $S \subseteq V$ . Then the following are equivalent:

- 1. S is dense.
- 2. For all  $x \in V$ , there exists a sequence  $(x_n)$  in S such that  $x_n \to x$ .
- 3. Every nonempty open subset of V intersects S.

#### Proof:

 $\langle 1 \rangle 1$ .  $1 \Leftrightarrow 2$ 

PROOF: Theorem 21.4.30.

 $\langle 1 \rangle 2$ .  $1 \Leftrightarrow 3$ 

Proof:

S is dense  $\Leftrightarrow$  the only closed set that includes S is V

 $\Leftrightarrow$  the only open set that does not intersect S is empty

**Definition 21.4.35** (Compact). Let V be a normed space and  $S \subseteq V$ . Then S is *compact* if and only if every sequence in S has a subsequence that converges to a limit in S.

Proposition 21.4.36. Every compact set is closed.

```
PROOF:  \langle 1 \rangle 1. \text{ Let: } V \text{ be a normed space.}   \langle 1 \rangle 2. \text{ Let: } C \subseteq V \text{ be compact.}   \langle 1 \rangle 3. \text{ Let: } (x_n) \text{ be a sequence in } C \text{ that converges to } l \in V.   \langle 1 \rangle 4. \text{ PICK a subsequence } (y_n) \text{ of } (x_n) \text{ that converges to } m \in C.   \langle 1 \rangle 5. \ l = m   \langle 1 \rangle 6. \ l \in C   \langle 1 \rangle 7. \text{ Q.E.D.}  PROOF: Proposition 21.4.28. 
  \Box
```

**Definition 21.4.37** (Bounded). Let V be a normed space and  $S \subseteq V$ . Then S is bounded iff there exists r > 0 such that  $S \subseteq B(0, r)$ .

**Proposition 21.4.38.** In  $\mathbb{R}^n$  and  $\mathbb{C}^n$ , the compact sets are the closed bounded sets.

# Proof:

- $\langle 1 \rangle 1$ . Let:  $C \subseteq K^n$
- $\langle 1 \rangle 2$ . If C is compact then C is closed.

Proof: Proposition 21.4.36.

- $\langle 1 \rangle 3$ . If C is compact then C is bounded.
  - $\langle 2 \rangle 1$ . Assume: C is compact.
  - $\langle 2 \rangle 2$ . Assume: for a contradiction C is not bounded.
  - $\langle 2 \rangle 3$ . For  $n \in \mathbb{N}$ , PICK  $x_n \in C$  with  $||x_n|| > n+1$ .
  - $\langle 2 \rangle 4$ . Pick a convergent subsequence  $(x_{n_r})$  that converges to  $l \in C$
  - $\langle 2 \rangle 5. \|x_{n_r}\| \to \|l\|$
  - $\langle 2 \rangle 6. \|x_{n_r}\| \to +\infty$
  - $\langle 2 \rangle$ 7. Q.E.D.

PROOF: This is a contradiction.

 $\langle 1 \rangle 4$ . If C is closed and bounded then C is compact.

PROOF: By the Bolzano-Weierstrass Theorem.

**Proposition 21.4.39.** Let V be a normed space and  $S \subseteq V$ . Then S is bounded if and only if, for every sequence  $(x_n)$  in S and every sequence  $(\lambda_n)$  in K, if  $\lambda_n \to 0$  then  $\|\lambda_n x_n\| \to 0$ .

# Proof:

- $\langle 1 \rangle 1$ . If S is bounded then, for every sequence  $(x_n)$  in S and every sequence  $(\lambda_n)$  in K, if  $\lambda_n \to 0$  then  $\|\lambda_n x_n\| \to 0$ .
  - $\langle 2 \rangle 1$ . Assume: S is bounded.
  - $\langle 2 \rangle 2$ . PICK r > 0 such that  $S \subseteq B(0, r)$ .
  - $\langle 2 \rangle 3$ . Let:  $(x_n)$  be a sequence in S.
  - $\langle 2 \rangle 4$ . Let:  $(\lambda_n)$  be a sequence in K.
  - $\langle 2 \rangle 5$ . Assume:  $\lambda_n \to 0$

- $\langle 2 \rangle 6$ . Let:  $\epsilon > 0$
- $\langle 2 \rangle$ 7. PICK N such that  $\forall n \geq N. |\lambda_n| < \epsilon/r$
- $\langle 2 \rangle 8. \ \forall n \geq N. \|\lambda_n x_n\| < \epsilon$
- $\langle 1 \rangle 2$ . If S is unbounded then there exists a sequence  $(x_n)$  in S and  $(\lambda_n)$  in K such that  $\lambda_n \to 0$  and  $||\lambda_n x_n|| \not\to 0$ .
  - $\langle 2 \rangle 1$ . S is unbounded.
  - $\langle 2 \rangle 2$ . For  $n \in \mathbb{N}$ , PICK  $x_n \in S$  such that  $||x_n|| > n$ .
  - $\langle 2 \rangle 3$ . For  $n \in \mathbb{N}$ , Let:  $\lambda_n = 1/n$  if n > 0, 1 if n = 0
  - $\langle 2 \rangle 4. \ \lambda_n \to 0$
- $\langle 2 \rangle 5$ .  $||\lambda_n x_n|| > 1$  for all n > 1

**Proposition 21.4.40.** In C([0,1]), the unit ball  $\overline{B(0,1)}$  is closed and bounded but not compact.

Proof:

 $\langle 1 \rangle 1$ .  $\overline{B(0,1)}$  is closed.

Proof: Proposition 21.4.23.

 $\langle 1 \rangle 2$ .  $\overline{B(0,\underline{1})}$  is bounded.

Proof:  $B(0,1) \subseteq B(0,2)$ .

- $\langle 1 \rangle 3$ .  $\overline{B(0,1)}$  is not compact.
  - $\langle 2 \rangle 1$ . For  $n \in \mathbb{N}$ ,

Let:  $x_n:[0,1]\to\mathbb{R}$  be the function  $x_n(t)=t^n$ .

- $\langle 2 \rangle 2$ . For  $n \in \mathbb{N}$ , we have  $x_n \in B(0,1)$ .
- $\langle 2 \rangle 3$ . No subsequence of  $(x_n)$  converges.

**Theorem 21.4.41** (Riesz's Lemma). Let X be a closed proper subspace of a normed space V. For every  $\epsilon \in (0,1)$ , there exists  $x_{\epsilon} \in V$  such that  $||x_{\epsilon}|| = 1$  and  $\forall x \in X. ||x_{\epsilon} - x|| \ge \epsilon$ .

Proof:

- $\langle 1 \rangle 1$ . Pick  $z \in E X$
- $\langle 1 \rangle 2$ . Let:  $d = \inf_{x \in X} ||z x||$
- $\langle 1 \rangle 3. \ d > 0$ 
  - $\langle 2 \rangle 1$ . Pick-  $\epsilon > 0$  such that  $B(z, \epsilon) \subseteq E X$
  - $\langle 2 \rangle 2. \ d \geq \epsilon$
- $\langle 1 \rangle 4$ . For all  $\epsilon \in (0,1)$ , choose  $y_{\epsilon} \in X$  such that

$$d \le ||z - y_{\epsilon}|| < d/\epsilon$$
.

 $\langle 1 \rangle$ 5. For  $\epsilon \in (0,1)$ , Let:

$$x_{\epsilon} = \frac{z - y_{\epsilon}}{\|z - y_{\epsilon}\|} .$$

 $\langle 1 \rangle 6$ . For all  $x \in X$  we have  $||x_{\epsilon} - x|| > \epsilon$ 

Proof:

$$||x_{\epsilon} - x|| = ||\frac{z - y_{\epsilon}}{||z - y_{\epsilon}||} - x||$$

$$= \frac{1}{||z - y_{\epsilon}||} ||z - y_{\epsilon} - ||z - y_{\epsilon}||x|| \qquad (y_{\epsilon} + ||z - y_{\epsilon}||x \in X)$$

$$\geq \frac{1}{||z - y_{\epsilon}||} d$$

$$> \epsilon$$

**Theorem 21.4.42.** Let V be a normed space. Then V is finite dimensional if and only if  $\overline{B(0,1)}$  is compact.

Proof:

- $\langle 1 \rangle 1$ . If V is finite dimensional then  $\overline{B(0,1)}$  is compact.
  - $\langle 2 \rangle 1$ . Assume: V is finite dimensional.
  - $\langle 2 \rangle 2$ . PICK a basis  $\{e_1, \ldots, e_n\}$ .
  - $\langle 2 \rangle 3$ . Assume: w.l.o.g.  $\forall \alpha_1, \dots, \alpha_n \in K . \|\alpha_1 e_1 + \dots + \alpha_n e_n\| = |\alpha_1| + \dots + |\alpha_n|$
  - $\langle 2 \rangle 4$ . Let:  $(x_m)$  be a sequence in  $\overline{B(0,1)}$
  - $\langle 2 \rangle 5$ . For  $m \in \mathbb{N}$ ,

Let:  $x_m = \alpha_{m1}e_1 + \cdots + \alpha_{mn}e_n$ 

- $\langle 2 \rangle 6$ . For  $m \in \mathbb{N}$  and  $i = 1, \ldots, n$ , we have  $|\alpha_{mi}| \leq 1$
- $\langle 2 \rangle$ 7. For i = 1, ..., n, PICK a convergent subsequence  $(\alpha_{m_r i})$  of  $(\alpha_{mi})$  in  $\mathbb{C}$  that converges to  $l_i$

PROOF: Since B(0,1) is compact in K.

 $\langle 2 \rangle 8$ .  $x_m$  converges to  $l_1 e_1 + \cdots + l_n e_n$ 

PROOF:

$$||x_m - l_1 e_1 - \dots - l_n e_n|| = ||(\alpha_{m1} - l_1)e_1 + \dots + (\alpha_{mn} - l_n)e_n||$$
  
=  $|\alpha_{m1} - l_1| + \dots + |\alpha_{mn} - l_n|$   
 $\to 0$  as  $m \to \infty$ 

- $\langle 1 \rangle 2$ . If V is infinite dimensional then B(0,1) is not compact.
  - $\langle 2 \rangle 1$ . Assume: V is infinite dimensional.
  - (2)2. Choose a sequence  $(x_n)$  such that  $||x_n|| = 1$  and  $||x_m x_n|| \ge 1/2$  for all  $m \ne n$ .
    - $\langle 3 \rangle 1$ . Assume: as induction hypothesis there exists a sequence  $(x_0, x_1, \dots, x_n)$  such that  $||x_i|| = 1$  and  $||x_i x_j|| \ge 1/2$  for  $i \ne j$
    - (3)2. PICK  $x_{n+1}$  such that  $||x_{n+1}|| = 1$  and  $||x_{n_1} x|| \ge 1/2$  for  $x \in \{x_1, \ldots, x_n\}$ .
  - $\langle 2 \rangle 3$ . Assume: for a contradiction  $(x_{n_r})$  is a subsequence that converges to l
  - $\langle 2 \rangle 4$ . For all r we have  $1/2 \le ||x_{n_r} l|| + ||x_{n_{r+1}} l||$  PROOF:

$$1/2 \le ||x_{n_r} - x_{n_{r+1}}||$$
  $(\langle 2 \rangle 2)$   $\le ||x_{n_r} - l|| + ||x_{n_{r+1}} - l||$  (Triangle Inequality)

 $\langle 2 \rangle$ 5. Q.E.D.

PROOF: This is a contradiction.

**Proposition 21.4.43.** Let U and V be normed spaces. Let  $f: U \to V$ . If f is continuous at one point, then it is continuous.

# Proof:

- $\langle 1 \rangle 1$ . Assume: f is continuous at  $x_0 \in U$ .
- $\langle 1 \rangle 2$ . Let:  $(x_n)$  be a sequence in U that converges to  $l \in U$ .
- $\langle 1 \rangle 3$ .  $x_n l + x_0 \to x_0$  as  $n \to \infty$
- $\langle 1 \rangle 4$ .  $f(x_n l + x_0) \to f(x_0)$  as  $n \to \infty$
- $\langle 1 \rangle 5. \ f(x_n) f(l) + f(x_0) \to f(x_0) \text{ as } n \to \infty$
- $\langle 1 \rangle 6. \ f(x_n) \to f(l) \text{ as } n \to \infty.$

**Definition 21.4.44** (Bounded). Let U and V be normed spaces and  $L: U \to V$  be a linear transformation. Then L is bounded iff there exists  $\alpha > 0$  such that  $\forall x \in U. ||L(x)|| \le \alpha ||x||$ .

**Theorem 21.4.45.** Let U and V be normed spaces. Let  $L: U \to V$  be a linear transformation. Then L is continuous if and only if it is bounded.

#### Proof:

- $\langle 1 \rangle 1$ . If L is continuous then L is bounded.
  - $\langle 2 \rangle 1$ . Assume: L is not bounded.
  - $\langle 2 \rangle 2$ . For  $n \in \mathbb{N}$ , choose  $x_n \in U$  such that  $||L(x_n)|| > (n+1)||x_n||$
  - $\langle 2 \rangle 3$ . For  $n \in \mathbb{N}$ ,

Let:  $y_n = x_n/(n+1)||x_n||$ 

- $\langle 2 \rangle 4$ .  $y_n \to 0$  as  $n \to \infty$
- $\langle 2 \rangle 5$ . For  $n \in \mathbb{N}$ , we have  $||L(y_n)|| > 1$
- $\langle 2 \rangle 6$ .  $L(y_n) \not\to 0 = L(0)$  as  $n \to \infty$
- $\langle 2 \rangle 7$ . L is not continuous.
- $\langle 1 \rangle 2$ . If L is bounded then L is continuous.
  - $\langle 2 \rangle$ 1. Let:  $\alpha > 0$  be such that  $\forall x \in U. ||L(x)|| \le \alpha ||x||$ Prove: L is continuous at 0.
  - $\langle 2 \rangle 2$ . Let:  $(x_n)$  be a sequence in U that converges to 0.
  - $\langle 2 \rangle 3$ .  $L(x_n) \to 0$  as  $n \to \infty$
  - $\langle 2 \rangle 4$ . Q.E.D.

Proof: Proposition 21.4.43.

**Corollary 21.4.45.1.** If U and V are finite dimensional normed spaces, then every linear transformation  $U \to V$  is continuous.

**Definition 21.4.46.** For U and V normed spaces, let  $\mathcal{B}(U,V)$  be the space of all bounded linear transformations  $U \to V$ . This is a subspace of the space of all linear transformations  $U \to V$ .

Define the uniform convergence norm on  $\mathcal{B}(U,V)$  by

$$||L|| = \sup\{||L(x)|| \mid x \in U, ||x|| = 1\}$$
.

We prove this is a norm.

 $\langle 2 \rangle 4. \ \|L + M\| \le \|L\| + \|M\|$ 

```
Proof:
```

```
\langle 1 \rangle 1. \ \forall L \in \mathcal{B}(U, V). ||L|| = 0 \Rightarrow L = 0
    \langle 2 \rangle 1. Let: L \in \mathcal{B}(U, V)
    \langle 2 \rangle 2. Assume: ||L|
    \langle 2 \rangle 3. For all x \in U, if ||x|| = 1 then ||L(x)|| = 0
    \langle 2 \rangle 4. Let: x \in U
              Prove: L(x) = 0
    \langle 2 \rangle5. Assume: w.l.o.g. x \neq 0
    \langle 2 \rangle 6. \ \|L(x/\|x\|)\| = 0
        Proof: \langle 2 \rangle 3
    \langle 2 \rangle 7. \ L(x/||x||) = 0
    \langle 2 \rangle 8. \ L(x) / ||x|| = 0
    \langle 2 \rangle 9. \ L(x) = 0
\langle 1 \rangle 2. \ \forall \lambda \in K. \forall L \in \mathcal{B}(U, V). \|\lambda L\| = |\lambda| \|L\|
    \langle 2 \rangle 1. Let: \lambda \in K
    \langle 2 \rangle 2. Let: L \in \mathcal{B}(U, V)
    \langle 2 \rangle 3. ||\lambda L|| = |\lambda| ||L||
        Proof:
                                                   \|\lambda L\| = \sup \|\lambda L(x)\|
                                                                 ||x|| = 1
                                                             = \sup (|\lambda| ||L(x)||)
                                                                 ||x|| = 1
                                                             = |\lambda| \sup ||L(x)||
                                                                      ||x|| = 1
                                                             = |\lambda| \|L\|
\langle 1 \rangle 3. The triangle inequality holds.
    \langle 2 \rangle 1. Let: L, M \in \mathcal{B}(U, V)
    \langle 2 \rangle 2. For all x \in U, if ||x|| = 1 then
                                           ||L(x) + M(x)|| \le ||L(x)|| + ||M(x)||.
    \langle 2 \rangle 3. For all x \in U, if ||x|| = 1 then
                                                ||L(x) + M(x)|| \le ||L|| + ||M||.
```

**Proposition 21.4.47.** Let U and V be normed spaces and  $L \in \mathcal{B}(U,V)$ . Then ||L|| is the least number such that  $\forall x \in U. ||L(x)|| \le ||L|| ||x|||$ .

**Theorem 21.4.48.** Let U and V be normed spaces. Let  $T: U \to V$  be a continuous linear transformation. Then the null space  $\mathcal{N}(T)$  is closed in U.

PROOF: If  $(x_n)$  is a sequence in  $\mathcal{N}(T)$  and  $x_n \to l$  then  $T(l) = \lim_{n \to \infty} T(x_n) = 0$  so  $l \in \mathcal{N}(T)$ .  $\square$ 

**Theorem 21.4.49** (Diagonal Theorem). Let V be a normed space. Let  $(x_{ij})_{i,j\in\mathbb{N}}$  be an infinite matrix in V. If:

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  - 1.  $\forall j \in \mathbb{N}. x_{ij} \to 0 \text{ as } i \to \infty;$
  - 2. Every strictly increasing sequence of natural numbers  $(p_i)$  has a subsequence  $(q_i)$  such that

$$\sum_{j=0}^{\infty} x_{q_i q_j} \to 0 \text{ as } i \to \infty$$

then  $x_{ii} \to 0$  as  $i \to \infty$ .

# Proof:

- $\langle 1 \rangle 1$ . Let: V be a normed space.
- $\langle 1 \rangle 2$ . Let:  $(x_{ij})_{i,j \in \mathbb{N}}$  be an infinite matrix in V.
- $\langle 1 \rangle 3$ . Assume:  $\forall j \in \mathbb{N}. x_{ij} \to 0 \text{ as } i \to \infty$
- $\langle 1 \rangle 4$ . Assume: Every strictly increasing sequence of natural numbers  $(p_i)$  has a subsequence  $(q_i)$  such that

$$\sum_{i=0}^{\infty} x_{q_i q_j} \to 0 \text{ as } i \to \infty$$

- $\langle 1 \rangle 5$ . Assume: for a contradiction  $x_{ii} \not\to 0$  as  $i \to \infty$
- $\langle 1 \rangle 6$ . PICK a strictly increasing sequence  $(p_i)$  and  $\epsilon > 0$  such that  $\forall i \in \mathbb{N}. ||x_{p_i p_i}|| \geq$ 
  - $\langle 2 \rangle 1$ . PICK  $\epsilon > 0$  such that, for all N, there exists  $n \geq N$  such that  $||x_{ii}|| \geq \epsilon$
  - $\langle 2 \rangle 2$ . Choose a sequence  $(p_i)$  such that, for all i, we have  $p_{i+1} \geq p_i + 1$  and  $||x_{p_ip_i}|| \geq \epsilon$
- $\langle 1 \rangle$ 7. PICK a subsequence  $(q_i)$  of  $(p_i)$  such that  $\sum_{j=0}^{\infty} x_{q_i q_j} \to 0$  as  $i \to \infty$ . Proof:  $\langle 1 \rangle 4$

 $\langle 1 \rangle$ 8. For all i we have  $x_{q_iq_j} \to 0$  as  $j \to \infty$ . PROOF: Since  $\sum_{j=0}^{\infty} x_{q_iq_j} < \infty$ .  $\langle 1 \rangle$ 9. For all j we have  $x_{q_iq_j} \to 0$  as  $i \to \infty$ .

Proof:  $\langle 1 \rangle 3$ 

- $\langle 1 \rangle 10$ . PICK a subsequence  $(r_i)$  of  $(q_i)$  such that, for all i, j with  $i \neq j$ , we have  $||x_{r_i r_j}|| < \epsilon/2^{j+2}$ .
  - $\langle 2 \rangle 1$ . Let:  $r_1 = q_1$
  - $\langle 2 \rangle 2$ . Assume: as induction hypothesis we have defined  $(r_1, \ldots, r_n)$  such that:

    - $r_1 < r_2 < \dots < r_n$  Whenever  $q_i > r_j$  then  $\|x_{q_ir_j}\| < \epsilon/2^{j+2}$
    - Whenever i < j then  $||x_{r_i r_j}|| < \epsilon/2^{j+2}$
  - $\langle 2 \rangle 3$ . Let:  $r_{n+1}$  be the first element of  $(q_i)$  such that  $r_{n+1} > r_n$ ,  $\|x_{q_i r_{n+1}}\| < \epsilon/2^{n+2}$  whenever  $q_i \geq r_{n+1}$ , and  $\|x_{r_j r_{n+1}}\| < \epsilon/2^{n+3}$  for j=1

PROOF: One exists by  $\langle 1 \rangle 8$ .

- $\langle 2 \rangle 4. \ r_1 < \dots < r_n < r_{n+1}$
- $\langle 2 \rangle 5$ . Whenever  $q_i > r_j$  then  $||x_{q_i r_j}|| < \epsilon/2^{j+2}$
- $\langle 2 \rangle 6$ . Whenever i < j then  $||x_{r_i r_j}|| < \epsilon/2^{j+2}$
- $\langle 1 \rangle 11$ . PICK a subsequence  $(s_i)$  of  $(r_i)$  such that  $\sum_{i=0}^{\infty} x_{s_i s_i} \to 0$  as  $i \to \infty$

PROOF: 
$$\langle 1 \rangle 4$$
  $\langle 1 \rangle 12$ .  $\forall i \in \mathbb{N}$ .  $\left\| \sum_{j=0}^{\infty} x_{s_i s_j} \right\| > \epsilon/2$  PROOF: 
$$\left\| \sum_{j=0}^{\infty} x_{s_i s_j} \right\| = \left\| x_{s_i s_i} + \sum_{j \neq i} x_{s_i s_j} \right\|$$
  $\geq \left\| \| x_{s_i s_i} \| - \left\| \sum_{j \neq i} x_{s_i s_j} \right\|$  (Proposition 21.4.5) 
$$= \| x_{s_i s_i} \| - \left\| \sum_{j \neq i} x_{s_i s_j} \right\| \qquad \left( \| x_{s_i s_i} \| > \epsilon, \left\| \sum_{j} x_{s_i s_j} \right\| \le \epsilon/2 \right)$$
  $\geq \| x_{s_i s_i} \| - \sum_{j \neq i} \| x_{s_i s_j} \| \qquad \text{(Triangle Inequality, )}$   $> \epsilon - \sum_{j \neq i} \epsilon/2^{j+2} \qquad (\langle 1 \rangle 6, \langle 1 \rangle 10)$   $> \epsilon/2$   $\langle 1 \rangle 13$ . Q.E.D. PROOF: This contradicts  $\langle 1 \rangle 11$ .

21.4.1 Functionals

**Definition 21.4.50** (Functional). Let K be either  $\mathbb{R}$  or  $\mathbb{C}$ . Let V be a normed space over K. A functional over V is a bounded linear transformation  $V \to K$ . The dual space of V is

$$V' := \mathcal{B}(V, K)$$
.

# 21.4.2 Contraction Mappings

**Definition 21.4.51** (Contraction). Let V be a normed space and  $A \subseteq V$ . Let  $f: A \to V$ . Then f is a *contraction* iff there exists a real number  $\alpha$  with  $0 < \alpha < 1$  such that, for all  $x, y \in A$ , we have

$$||f(x) - f(y)|| \le \alpha ||x - y||$$
.

Proposition 21.4.52. Every contraction mapping is continuous.

# 21.5 Banach Spaces

**Definition 21.5.1** (Cauchy Sequence). A sequence  $(x_n)$  in a normed space is a *Cauchy sequence* iff, for every  $\epsilon > 0$ , there exists N such that  $\forall m, n \geq N . ||x_m - x_n|| < \epsilon$ .

**Theorem 21.5.2.** Let V be a normed space. Let  $(x_n)$  be a sequence in V. Then the following are equivalent.

- 1.  $(x_n)$  is Cauchy.
- 2. For every pair of strictly increasing sequences of natural numbers  $(p_n)$  and  $(q_n)$ , we have  $||x_{p_n} x_{q_n}|| \to 0$  as  $n \to \infty$ .
- 3. For every strictly increasing sequence of natural numbers  $(p_n)$ , we have  $||x_{p_{n+1}} x_{p_n}|| \to 0$  as  $n \to \infty$ .

# Proof:

- $\langle 1 \rangle 1. \ 1 \Rightarrow 2$ 
  - $\langle 2 \rangle 1$ . Assume:  $(x_n)$  is Cauchy.
  - $\langle 2 \rangle 2$ . Let:  $(p_n)$  and  $(q_n)$  be a pair of increasing sequences of natural numbers.
  - $\langle 2 \rangle 3$ . Let:  $\epsilon > 0$
  - $\langle 2 \rangle 4$ . PICK N such that  $\forall m, n \geq N . ||x_m x_n|| < \epsilon$
  - $\langle 2 \rangle 5. \ \forall n \geq N.p_n, q_n \geq N$
  - $\langle 2 \rangle 6. \ \forall n \geq N. \|x_{p_n} x_{q_n}\| < \epsilon$
- $\langle 1 \rangle 2. \ 2 \Rightarrow 3$

PROOF: Trivial.

- $\langle 1 \rangle 3. \ 2 \Rightarrow 1$ 
  - $\langle 2 \rangle 1$ . Assume:  $(x_n)$  is not Cauchy.
  - $\langle 2 \rangle 2$ . PICK  $\epsilon > 0$  such that, for all N, there exist  $m, n \geq N$  such that  $||x_m x_n|| \geq \epsilon$
  - $\langle 2 \rangle$ 3. Pick strictly increasing sequences of natural numbers  $(p_n)$  and  $(q_n)$  such that, for all n,  $||x_{p_n} x_{q_n}|| \ge \epsilon$ 
    - $\langle 3 \rangle 1$ . Assume: as induction hypothesis we have chosen  $(p_1, \ldots, p_n)$  and  $(q_1, \ldots, q_n)$  strictly increasing such that  $\forall i. ||x_{p_i} x_{q_i}|| \ge \epsilon$
  - $\langle 3 \rangle 2$ . Pick  $p_{n+1}, q_{n+1} \geq \max(p_n, q_n)$  such that  $||x_{p_{n+1}} x_{q_{n+1}}|| \geq \epsilon$
  - $\langle 2 \rangle 4$ . 2 is false.
- $\langle 1 \rangle 4. \ 3 \Rightarrow 2$ 
  - $\langle 2 \rangle 1$ . Assume:  $(p_n)$  and  $(q_n)$  are strictly increasonig sequences such that  $\|x_{p_n} x_{q_n}\| \not\to 0$  as  $n \to \infty$ .
  - $\langle 2 \rangle$ 2. PICK  $\epsilon > 0$  such that, for all N, there exists  $n \geq N$  such that  $\|x_{p_n} x_{q_n}\| \geq \epsilon$
  - $\langle 2 \rangle$ 3. Choose a strictly increasing sequence  $(r_n)$  such that, for all n, we have  $||x_{r_{2n}} x_{r_{2n+1}}|| \ge \epsilon$ 
    - $\langle 3 \rangle 1$ . Assume: as induction hypothesis we have chosen  $(r_0, r_1, \ldots, r_{2n+1})$  such that, for  $i = 0, 1, \ldots, n$ , we have  $||x_{r_{2i}} x_{r_{2i+1}}|| \ge \epsilon$
    - $\langle 3 \rangle 2$ . Pick  $i, j \geq r_{2n+1}$  such that  $||x_i x_j|| \geq \epsilon$
    - $\langle 3 \rangle 3$ . Set  $r_{2n+2} = \min(i, j)$  and  $r_{2n+3} = \max(i, j)$
- $\langle 2 \rangle 4$ . 3 is false.

**Proposition 21.5.3.** Every convergent sequence is a Cauchy sequence.

Proof:

- $\langle 1 \rangle 1$ . Let:  $(x_n)$  be a convergent sequence in a normed space V with limit l.
- $\langle 1 \rangle 2$ . Let:  $\epsilon > 0$
- $\langle 1 \rangle 3$ . PICK N such that  $\forall n \geq N . ||x_n l|| < \epsilon/2$
- $\langle 1 \rangle 4. \ \forall m, n \geq N. ||x_m x_n|| < \epsilon$

**Proposition 21.5.4.** Let  $\mathcal{P}([0,1])$  be the space of polynomials on [0,1] under the norm of uniform convergence. For  $n \in \mathbb{N}$ , let  $P_n = 1 + x + x^2/2! + \cdots + x^n/n!$ . Then  $(P_n)$  is Cauchy but does not converge, since  $e^x$  is not a polynomial.

Proof: Easy.

**Proposition 21.5.5.** If  $(x_n)$  is a Cauchy sequence in a normed space V, then  $(\|x_n\|)$  converges in  $\mathbb{R}$ .

Proof:

- $\langle 1 \rangle 1$ . ( $||x_n||$ ) is Cauchy.
  - $\langle 2 \rangle 1$ . Let:  $\epsilon > 0$
  - $\langle 2 \rangle 2$ . PICK N such that  $\forall m, n \geq N . ||x_m x_n|| < \epsilon$
  - $\langle 2 \rangle 3. \ \forall m, n \geq N. |||x_m|| ||x_n||| < \epsilon$

Proof: Proposition 21.4.5.

 $\langle 1 \rangle 2$ . Q.E.D.

PROOF: Since  $\mathbb{R}$  is complete.

Corollary 21.5.5.1. Every Cauchy sequence is bounded.

**Definition 21.5.6** (Banach space). A normed space is *complete* or a *Banach* space iff every Cauchy sequence converges.

**Proposition 21.5.7.** For all  $p \ge 1$ , the space  $l^p$  in complete.

Proof:

- $\langle 1 \rangle 1$ . Let:  $(a_n)$  be a Cauchy sequence in  $l^p$ .
- $\langle 1 \rangle 2$ . For  $n \in \mathbb{N}$ ,

Let:  $a_n = (\alpha_{n0}, \alpha_{n1}, \ldots)$ 

 $\langle 1 \rangle 3$ . For all  $\epsilon > 0$ , there exists N such that  $\forall m, n \geq N$ 

$$\sum_{k=0}^{\infty} |\alpha_{mk} - \alpha_{nk}|^p < \epsilon .$$

Proof:  $\langle 1 \rangle 1$ 

 $\langle 1 \rangle 4$ . For all  $k \in \mathbb{N}$  and  $\epsilon > 0$ , there exists N such that  $\forall m, n \geq N$ 

 $|\alpha_{mk} - \alpha_{nk}| < \epsilon$ .

Proof:  $\langle 1 \rangle 3$ 

 $\langle 1 \rangle$ 5. For all  $k \in \mathbb{N}$ , the sequence  $(\alpha_{nk})_n$  converges in  $\mathbb{C}$ .

PROOF:  $\mathbb{C}$  is complete.

 $\langle 1 \rangle 6$ . For  $k \in \mathbb{N}$ ,

Let:

$$\alpha_k = \lim_{n \to \infty} \alpha_{nk} .$$

- $\langle 1 \rangle 7$ . Let:  $a = (\alpha_k)_k$
- $\langle 1 \rangle 8$ . For all  $\epsilon > 0$ , there exists N such that  $\forall n \geq N$

$$\sum_{k=0}^{\infty} |\alpha_k - \alpha_{nk}|^p < \epsilon .$$

PROOF: Take the limit  $m \to \infty$  in  $\langle 1 \rangle 4$ .

- $\langle 1 \rangle 9. \ a \in l^p$ 
  - $\langle 2 \rangle 1$ . PICK N such that  $\forall n \geq N$ .  $\sum_{k=0}^{\infty} |\alpha_k \alpha_{nk}|^p < 1$
  - $\langle 2 \rangle 2$ .  $a a_N \in l^p$
  - $\langle 2 \rangle 3$ . Q.E.D.

PROOF: Since  $l^p$  is closed under +.

 $\langle 1 \rangle 10$ .  $a_n \to a$  as  $n \to \infty$ .

PROOF: Immediate from  $\langle 1 \rangle 8$ .

**Proposition 21.5.8.** For any real number a, b with a < b, the space C([a,b]) is complete.

# Proof:

- $\langle 1 \rangle 1$ . Let:  $(f_n)$  be a Cauchy sequence in  $\mathcal{C}([a,b])$
- $\langle 1 \rangle 2$ . For all  $\epsilon > 0$ , there exists N such that, for all  $m, n \geq N$  and  $x \in [a, b]$ ,

$$|f_m(x) - f_n(x)| < \epsilon$$
.

- $\langle 1 \rangle 3$ . For all  $x \in [a, b]$ ,  $(f_n(x))_n$  is Cauchy.
- $\langle 1 \rangle 4$ . Let:  $f : [a, b] \to \mathbb{C}$  be the function

$$f(x) = \lim_{n \to \infty} f_n(x) .$$

 $\langle 1 \rangle 5$ . For all  $\epsilon > 0$ , there exists N such that, for all  $n \geq N$  and  $x \in [a, b]$ ,

$$|f_n(x) - f(x)| < \epsilon$$
.

PROOF: Take the limit  $m \to \infty$  in  $\langle 1 \rangle 2$ .

- $\langle 1 \rangle 6$ . f is continuous.
  - $\langle 2 \rangle 1$ . Let:  $x_0 \in [a, b]$
  - $\langle 2 \rangle 2$ . Let:  $\epsilon > 0$
  - $\langle 2 \rangle 3$ . PICK N such that, for all  $n \geq N$  and  $y \in [a, b]$ , we have

$$|f_n(y) - f(y)| < \epsilon/3$$
.

Proof:  $\langle 1 \rangle 5$ 

- $\langle 2 \rangle 4$ . PICK  $\delta > 0$  such that, for all  $y \in [a, b]$  with  $|x_0 y| < \delta$ , we have  $|f_N(x_0) f_N(y)| < \epsilon/3$
- $\langle 2 \rangle$ 5. For all  $y \in [a, b]$ , if  $|x_0 y| < \delta$  then  $|f(x_0) f(y)| < \epsilon$ PROOF:

$$|f(x_0) - f(y)|$$

$$\leq |f(x_0) - f_N(x_0)| + |f_N(x_0) - f_N(y)| + |f_N(y) - f(y)| \quad \text{(Triangle Inequality)}$$
$$< \epsilon/3 + \epsilon/3 + \epsilon/3 \qquad (\langle 2 \rangle 3, \, \langle 2 \rangle 4)$$

 $=\epsilon$ 

 $\langle 1 \rangle 7$ .  $f_n \to f$  as  $n \to \infty$ .

PROOF: Immediate from  $\langle 1 \rangle 5$ .

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**Definition 21.5.9** (Convergent Series). Let  $(x_n)$  be a sequence in a normed space V. We say the series  $\sum_{n=0}^{\infty} x_n$  is convergent iff the sequence  $(\sum_{n=0}^{N} x_n)_N$  is convergent, and then we write  $\sum_{n=0}^{\infty} x_n = l$  for  $\lim_{N\to\infty} \sum_{n=0}^{N} x_n = l$ .

**Definition 21.5.10** (Absolutely Convergent Series). Let  $(x_n)$  be a sequence in a normed space V. We say the series  $\sum_{n=0}^{\infty} x_n$  is absolutely convergent iff the series  $\sum_{n=0}^{\infty} ||x_n||$  converges in  $\mathbb{R}$ .

**Theorem 21.5.11.** A normed space is complete if and only if every absolutely convergent series is convergent.

# Proof:

- $\langle 1 \rangle 1$ . Let: V be a normed space.
- $\langle 1 \rangle 2$ . If V is complete then every absolutely convergent series is convergent.
  - $\langle 2 \rangle 1$ . Assume: V is complete.
  - $\langle 2 \rangle 2$ . Let:  $(x_n)$  be a sequence in V.
  - $\langle 2 \rangle$ 3. Assume:  $\sum_{n=0}^{\infty} \|x_n\|$  converges. Prove:  $(\sum_{n=0}^{N} x_n)_N$  is Cauchy.
  - $\langle 2 \rangle 4$ . Let:  $\epsilon > 0$
  - $\langle 2 \rangle$ 5. PICK N such that  $\sum_{n=N}^{\infty} ||x_n|| < \epsilon$

  - $\langle 2 \rangle$ 6. Let:  $m > n \ge N$   $\langle 2 \rangle$ 7.  $\| \sum_{k=0}^{m} x_k \sum_{k=0}^{n} x_k \| < \epsilon$

$$\| \sum_{k=0}^{m} x_k - \sum_{k=0}^{n} x_k \| = \| \sum_{k=n+1}^{m} x_k \|$$

$$\leq \sum_{k=n+1}^{m} \| x_k \|$$

$$\leq \sum_{k=n+1}^{\infty} \| x_k \|$$

$$\leq \epsilon$$

- $\langle 1 \rangle 3$ . If every absolutely convergent series is convergent then V is complete.
  - $\langle 2 \rangle 1$ . Assume: Every absolutely convergent series is convergent.
  - $\langle 2 \rangle 2$ . Let:  $(x_n)$  be a Cauchy sequence in V.
  - $\langle 2 \rangle 3$ . Choose an increasing sequence of natural numbers  $(p_k)$  such that, for all  $m, n \geq p_k$ , we have

$$||x_m - x_n|| < 2^{-k}$$
.

- $||x_m x_n|| < 2^{-k} .$   $\langle 2 \rangle 4$ .  $\sum_{k=0}^{\infty} ||x_{p_{k+1}} x_{p_k}||$  is absolutely convergent.  $\langle 2 \rangle 5$ .  $\sum_{k=0}^{\infty} ||x_{p_{k+1}} x_{p_k}||$  is convergent.
- $\langle 2 \rangle 6$ .  $(x_{p_k})$  converges.
- $\langle 2 \rangle 7$ . Let:  $l = \lim_{k \to \infty} x_{p_k}$
- $\langle 2 \rangle 8. \ x_n \to l \text{ as } n \to \infty.$

Proof:

$$||x_n - l|| \le ||x_n - x_{p_n}|| + ||x_{p_n} - l||$$
  
  $\to 0$  as  $n \to \infty$  (Theorem 21.5.2)

**Proposition 21.5.12.** A closed subspace of a Banach space is a Banach space.

#### Proof:

- $\langle 1 \rangle 1$ . Let: V be a Banach space.
- $\langle 1 \rangle 2$ . Let: U be a closed subspace of V.
- $\langle 1 \rangle 3$ . Let:  $(x_n)$  be a Cauchy sequence in U.
- $\langle 1 \rangle 4$ .  $(x_n)$  is a Cauchy sequence in V.
- $\langle 1 \rangle 5$ . Let: *l* be the limit of  $(x_n)$  in *V*.
- (1)6. Let t be the finite of  $(x_n)$

 $\langle 1 \rangle 6. \ l \in U$ 

Proof: Proposition 21.4.28.

**Definition 21.5.13** (Completion). Let V be a normed space. A *completion* of V consists of a Banach space W and a function  $\phi: V \to W$  such that:

- $\phi$  is injective.
- $\phi$  is a linear transformation.
- $\phi$  preserves the norm.
- $\phi(V)$  is dense in W.

**Definition 21.5.14** (Equivalent Cauchy Sequences). Let V be a normed space. Let  $(x_n)$  and  $(y_n)$  be Cauchy sequences in V. Then  $(x_n)$  and  $(y_n)$  are equivalent,  $(x_n) \sim (y_n)$ , iff  $||x_n - y_n|| \to 0$  as  $n \to \infty$ .

**Proposition 21.5.15.** Equivalence is an equivalence relation on the set of Cauchy sequences.

**Proposition 21.5.16.** If  $(x_n) \sim (y_n)$  then  $\lim_{n\to\infty} ||x_n|| = \lim_{n\to\infty} ||y_n||$ .

**Theorem 21.5.17.** Let V be a normed space. Let W be the quotient set of all Cauchy sequences modulo  $\sim$ . Define +,  $\cdot$  and  $\| \cdot \|$  on W by

$$[(x_n)] + [(y_n)] = [(x_n + y_n)]$$
$$\lambda[(x_n)] = [(\lambda x_n)]$$
$$\|[(x_n)]\| = \lim_{n \to \infty} \|x_n\|$$

Define  $\phi: V \to W$  by  $\phi(x)$  is the constant sequence (x). Then  $\phi: V \to W$  is the completion of V.

# Proof:

- $\langle 1 \rangle 1$ . +, · and  $\| \|$  are well defined.
- $\langle 1 \rangle 2$ . W is a normed space.
- $\langle 1 \rangle 3$ .  $\phi(V)$  is dense in W.

PROOF: For any  $[(x_n)] \in W$  we have  $[(x_n)] = \lim_{n \to \infty} \phi(x_n)$ .

 $\langle 1 \rangle 4$ . W is complete.

- $\langle 2 \rangle 1$ . Let:  $(X_n)$  be a Cauchy sequence in W.
- $\langle 2 \rangle 2$ . For all  $n \in \mathbb{N}$ , PICK  $x_n \in V$  such that  $\|\phi(x_n) X_n\| < 1/n$ PROOF:  $\langle 1 \rangle 3$
- $\langle 2 \rangle$ 3. For all m, n we have  $||x_n x_m|| \le ||X_n X_m|| + 1/n + 1/m$ PROOF:

$$||x_n - x_m|| = ||\phi(x_n) - \phi(x_m)||$$

$$\leq ||\phi(x_n) - X_n|| + ||X_n - X_m|| + ||\phi(x_m) - X_m||$$

$$\leq ||X_n - X_m|| + 1/n + 1/m$$

- $\langle 2 \rangle 4$ .  $(x_n)$  is a Cauchy sequence in V.
- $\langle 2 \rangle$ 5. Let:  $X = [(x_n)]$ Prove:  $X_n \to X$  as  $n \to \infty$
- $\langle 2 \rangle 6$ .  $||X_n X|| \to 0$  as  $n \to \infty$

Proof:

$$||X_n - X|| \le ||X_n - \phi(x_n)|| + ||\phi(x_n) - X||$$

$$< ||\phi(x_n) - X|| + 1/n$$

$$\to 0$$

- $\langle 1 \rangle 5$ .  $\phi$  is injective.
- $\langle 1 \rangle 6$ .  $\phi$  is a linear transformation.
- $\langle 1 \rangle 7$ .  $\phi$  preserves the norm.

**Theorem 21.5.18.** Let U be a normed space and V a Banach space. Then  $\mathcal{B}(U,V)$  is a Banach space.

# Proof:

- $\langle 1 \rangle 1$ . Let:  $(L_n)$  be a Cauchy sequence in  $\mathcal{B}(U,V)$
- $\langle 1 \rangle 2$ . For all  $x \in U$ , we have  $(L_n(x))$  is a Cauchy sequence in V.
  - $\langle 2 \rangle 1$ . Let:  $x \in U$
  - $\langle 2 \rangle 2$ . Assume: w.l.o.g.  $x \neq 0$
  - $\langle 2 \rangle 3$ . Let:  $\epsilon > 0$
  - $\langle 2 \rangle 4$ . PICK N such that  $\forall m, n \geq N . ||L_m L_n|| < \epsilon / ||x||$
  - $\langle 2 \rangle 5. \ \forall m, n \ge N. ||L_m(x) L_n(x)|| < \epsilon$

PROOF:  $||L_m(x) - L_n(x)|| \le ||L_m - L_n|| ||x|| < \epsilon$ 

- $\langle 1 \rangle 3$ . Define  $L: U \to V$  by  $L(x) = \lim_{n \to \infty} L_n(x)$
- $\langle 1 \rangle 4. \ L \in \mathcal{B}(U,V)$ 
  - $\langle 2 \rangle 1$ . L is linear.
    - $\langle 3 \rangle 1$ . Let:  $\lambda, \mu \in K$  and  $x, y \in U$
    - $\langle 3 \rangle 2$ .  $L(\lambda x + \mu y) = \lambda L(x) + \mu L(y)$

PROOF:

$$L(\lambda x + \mu y) = \lim_{n \to \infty} L_n(\lambda x + \mu y)$$
$$= \lim_{n \to \infty} (\lambda L_n(x) + \mu L_n(y))$$
$$= \lambda L(x) + \mu L(y)$$

- $\langle 2 \rangle 2$ . L is bounded.
  - $\langle 3 \rangle 1$ . PICK N such that  $\forall m, n \geq N . ||L_m L_n|| < 1$

```
PROVE: \forall x \in U. ||L(x)|| \le (||L_N|| + 1) ||x||
        \langle 3 \rangle 2. \ \forall n \geq N. ||L_n|| \leq ||L_N|| + 1
           PROOF: Since |||L_n|| - ||L_N||| \le ||L_n - L_N|| < 1.
        \langle 3 \rangle 3. Let: x \in U
        \langle 3 \rangle 4. \ \|L(x)\| \le (\|L_N\| + 1)\|x\|
           Proof:
                                                   ||L(x)|| = \left\|\lim_{n \to \infty} L_n(x)\right\|
                                                               = \lim_{n \to \infty} \|L_n(x)\|
                                                                \leq \lim_{n\to\infty} ||L_n|| ||x||
                                                                \leq (\|L_N\| + 1)\|x\|
\langle 1 \rangle 5. L_n \to L as n \to \infty
    \langle 2 \rangle 1. Let: \epsilon > 0
    \langle 2 \rangle 2. PICK N such that \forall m, n \geq N . ||L_m - L_n|| < \epsilon/4
    \langle 2 \rangle 3. Let: n \geq N
    \langle 2 \rangle 4. For all x \in U we have ||L_n(x) - L(x)|| < (\epsilon/2)||x||
        \langle 3 \rangle 1. Let: x \in U
        \langle 3 \rangle 2. For all m \geq N we have ||L_n(x) - L_m(x)|| < (\epsilon/4)||x||
        \langle 3 \rangle 3. \| L_n(x) - L(x) \| \le (\epsilon/4) \| x \|
           PROOF: Taking the limit as m \to \infty.
    \langle 2 \rangle 5. \ \|L_n - L\| \le \epsilon/2
    \langle 2 \rangle 6. ||L_n - L|| < \epsilon
П
```

Corollary 21.5.18.1. The dual space of a normed space is a Banach space.

**Theorem 21.5.19.** Let U be a normed space and V a Banach space. Let W be a subspace of U. Let  $L: W \to V$  be a bounded linear transformation. Then L has a unique extension to a bounded linear transformation  $\overline{W} \to V$ .

### Proof:

- $\langle 1 \rangle 1$ . Define  $L' : \overline{W} \to V$  as follows. Given  $x \in \overline{W}$ , pick a sequence  $(x_n)$  in W that converges to x. Then  $L'(x) = \lim_{n \to \infty} L(x_n)$ 
  - $\langle 2 \rangle 1$ . For all  $x \in \overline{W}$ , there exists a sequence  $(x_n)$  in W that converges to x. PROOF: Theorem 21.4.30.
  - $\langle 2 \rangle 2$ . For any sequence  $(x_n)$  in W that converges in  $\overline{W}$ , we have  $(L(x_n))$  converges in V.
    - $\langle 3 \rangle 1$ . Let:  $(x_n)$  be a sequence in W that converges in  $\overline{W}$
    - $\langle 3 \rangle 2$ .  $(x_n)$  is Cauchy.
    - $\langle 3 \rangle 3$ .  $(L(x_n))$  is Cauchy.

PROOF: For any strictly increasing sequence of natural numbers  $(p_n)$ , we have  $||L(x_{p_{n+1}}) - L(x_{p_n})|| \le ||L|| ||x_{p_{n+1}} - x_{p_n}|| \to 0$  as  $n \to \infty$ .

 $\langle 3 \rangle 4$ . Q.E.D.

PROOF: W is a Banach space.

 $\langle 2 \rangle$ 3. If  $(x_n)$  and  $(y_n)$  are sequences in W that converge to the same point in  $\overline{W}$ , then  $\lim_{n\to\infty} L(x_n) = \lim_{n\to\infty} L(y_n)$ 

PROOF: Since  $||L(x_n) - L(y_n)|| \le ||L|| ||x_n - y_n|| \to 0$ .

 $\langle 1 \rangle 2$ . L' extends L

PROOF: For  $x \in W$  we have the constant sequence (x) converges to x, and the constant sequence (L(x)) converges to L(x), so L'(x) = L(x).

- $\langle 1 \rangle 3$ . L' is a linear transformation.
  - $\langle 2 \rangle 1$ . Let:  $\lambda, \mu \in K$  and  $x, y \in \overline{W}$
  - $\langle 2 \rangle 2$ . PICK sequences  $(x_n)$ ,  $(y_n)$  in W that converge to x and y respectively.
  - $\langle 2 \rangle 3. \ \lambda x_n + \mu y_n \to \lambda x + \mu y$
  - $\langle 2 \rangle 4$ . Q.E.D.

Proof:

$$L'(\lambda x + \mu y) = \lim_{n \to \infty} L(\lambda x_n + \mu y_n)$$
$$= \lim_{n \to \infty} (\lambda L(x_n) + \mu L(y_n))$$
$$= \lambda L'(x) + \mu L'(y)$$

- $\langle 1 \rangle 4$ . L' is bounded.
  - $\langle 2 \rangle 1$ . Let:  $x \in \overline{W}$
  - $\langle 2 \rangle 2$ . PICK a sequence  $(x_n)$  in W that converges to x.
  - $\langle 2 \rangle 3. \| L'(x) \| \le \| L \| \| x \|$

Proof:

$$||L'(x)|| = ||\lim_{n \to \infty} L(x_n)||$$

$$= \lim_{n \to \infty} ||L(x_n)||$$

$$\leq ||L|| \lim_{n \to \infty} ||x_n||$$

$$= ||L|||x||$$

- $\langle 1 \rangle 5.$  If  $L'': \overline{W} \to V$  is a bounded linear transformation that extends L, then L'' = L'.
  - $\langle 2 \rangle 1$ . Let:  $x \in \overline{W}$
  - $\langle 2 \rangle 2$ . PICK a sequence  $(x_n)$  in W that converges to x.
  - $\langle 2 \rangle 3. \ L''(x) = L'(x)$

Proof:

$$L''(x) = \lim_{n \to \infty} L''(x_n)$$
$$= \lim_{n \to \infty} L(x_n)$$
$$= L'(x_n)$$

**Theorem 21.5.20** (Banach-Steinhaus). Let X be a Banach space and Y a normed space. Let  $\mathcal{T}$  be a set of bounded linear transformations from X into Y. Assume that, for all  $x \in X$ , there exists  $M_x > 0$  such that, for all  $T \in \mathcal{T}$ , we have  $||T(x)|| \leq M_x$ . Then there exists M > 0 such that, for all  $T \in \mathcal{T}$ , we have  $||T|| \leq M$ .

# Proof:

- $\langle 1 \rangle 1$ . Let: X be a Banach space.
- $\langle 1 \rangle 2$ . Let: Y be a normed space.

- $\langle 1 \rangle 3$ . Let:  $\mathcal{T}$  be a set of bounded linear transformations.
- $\langle 1 \rangle 4$ . Assume: For all  $x \in X$ , there exists  $M_x > 0$  such that, for all  $T \in \mathcal{T}$ , we have  $||T(x)|| \leq M_x$ .
- $\langle 1 \rangle$ 5. Assume: for a contradiction there is no M > 0 such that, for all  $T \in \mathcal{T}$ , we have  $||T|| \leq M$
- $\langle 1 \rangle 6$ . For every positive integer n, choose  $T_n \in \mathcal{T}$  such that  $||T_n|| > n2^n$ .
- $\langle 1 \rangle$ 7. For every positive integer n, choose  $x_n \in U$  such that  $||x_n|| = 1$  and  $||T_n(x_n)|| > n2^n.$
- $\langle 1 \rangle 8$ . For every positive integer n we have

$$\left\| \frac{1}{n} T_n \left( \frac{x_n}{2^n} \right) \right\| > 1 .$$

 $\langle 1 \rangle 9$ . For positive integers i and j Let:

$$y_{ij} = \frac{1}{i} T_i(\frac{x_j}{2^j})$$

 $\begin{array}{ll} \langle 1 \rangle 10. \ \ \mathrm{Let:} \ z = \sum_{j=1}^{\infty} \frac{x_j}{2^j} \\ \langle 2 \rangle 1. \ \sum_{j=1}^{\infty} \frac{x_j}{2^j} \ \ \mathrm{is \ absolutely \ convergent.} \\ \mathrm{Proof:} \end{array}$ 

$$\sum_{j=1}^{\infty} \left\| \frac{x_j}{2^j} \right\| = \sum_{j=1}^{\infty} \frac{1}{2^j} \tag{\langle 1 \rangle 7}$$

 $\langle 2 \rangle 2$ . Q.E.D.

Proof: Theorem 21.5.11.

- $\langle 1 \rangle 11$ . PICK C > 0 such that, for all i, we have  $\left\| \sum_{j=1}^{\infty} y_{ij} \right\| \leq C/i$ .
  - $\langle 2 \rangle 1$ . PICK C > 0 such that, for all  $T \in \mathcal{T}$ , we have  $||T(z)|| \leq C$ . Proof:  $\langle 1 \rangle 4$ .
  - $\langle 2 \rangle 2$ . For all i we have  $\left\| \sum_{j=1}^{\infty} y_{ij} \right\| \leq C/i$

TROOF. 
$$\left\| \sum_{j=1}^{\infty} y_{ij} \right\| = \left\| \sum_{j=1}^{\infty} \frac{1}{i} T_i \left( \frac{x_j}{2^j} \right) \right\|$$
 (\$\langle 1 \) 
$$= \frac{1}{i} \left\| T_i \left( \sum_{j=1}^{\infty} \frac{x_j}{2^j} \right) \right\|$$
 (\$\tau \text{continuous by } \langle 1 \rangle 3\$) 
$$= \frac{1}{i} \| T_i(z) \|$$
 (\$\langle 1 \rangle 1 \rangle 1\$) 
$$\leq \frac{C}{i}$$
 (\$\langle 2 \rangle 1\$) Is For any increasing sequence of positive integers \$(a\_i)\$, we have \$\sum\_{i=1}^{\infty}\$.

- $\langle 1 \rangle 13$ . For any increasing sequence of positive integers  $(q_i)$ , we have  $\sum_{j=0}^{\infty} y_{q_i q_j} \to$ 0 as  $i \to \infty$

PROOF: Similar.

 $\langle 1 \rangle 14$ . For all j we have  $y_{ij} \to 0$  as  $i \to \infty$ 

PROOF: From  $\langle 1 \rangle 9$  and the fact that  $T_i$  is continuous.

 $\langle 1 \rangle 15$ .  $y_{ii} \to 0$  as  $i \to \infty$ .

PROOF: By the Diagonal Theorem.

 $\langle 1 \rangle 16$ . Q.E.D.

PROOF: This contradicts  $\langle 1 \rangle 8$ .

**Theorem 21.5.21** (Banach Fixed Point Theorem). Let V be a Banach space. Let  $F \subseteq V$  be closed and nonempty. Let  $f: F \to F$  be a contraction. Then there exists a unique  $z \in F$  such that f(z) = z.

## Proof:

- $\langle 1 \rangle 1$ . PICK  $\alpha$  such that  $0 < \alpha < 1$  and  $\forall x, y \in V || f(x) f(y) || \leq \alpha || x y ||$ .
- $\langle 1 \rangle 2$ . Pick  $x_0 \in F$
- $\langle 1 \rangle 3$ . Extend to the sequence  $(x_n)$  in F by defining  $x_{n+1} := f(x_n)$ .
- $\langle 1 \rangle 4$ .  $(x_n)$  is Cauchy.
  - $\langle 2 \rangle 1. \ \forall n \in \mathbb{N}. \|x_{n+1} x_n\| \le \alpha^n \|x_1 x_0\|$
  - $\langle 2 \rangle 2$ . Let:  $\epsilon > 0$
  - $\langle 2 \rangle$ 3. PICK N such that  $\frac{\|x_1 x_0\|}{1 \alpha} \alpha^N \le \epsilon$   $\langle 2 \rangle$ 4.  $\forall m, n \ge N. \|x_n x_m\| < \epsilon$
  - - $\langle 3 \rangle 1$ . Let:  $m, n \geq N$
    - $\langle 3 \rangle 2$ . Assume: w.l.o.g. m < n
    - $\langle 3 \rangle 3. \|x_n x_m\| < \epsilon$

Proof:

$$||x_{n} - x_{m}|| \leq ||x_{n} - x_{n-1}|| + ||x_{n-1} - x_{n-2}|| + \dots + ||x_{m+1} - x_{m}||$$

$$\leq (\alpha^{n-1} + \alpha^{n-2} + \dots + \alpha^{m})||x_{1} - x_{0}||$$

$$< \frac{||x_{1} - x_{0}||}{1 - \alpha} \alpha^{m}$$

$$\leq \frac{||x_{1} - x_{0}||}{1 - \alpha} \alpha^{N}$$

$$\leq \epsilon$$

- $\langle 1 \rangle 5$ . Let:  $z = \lim_{n \to \infty} x_n$
- $\langle 1 \rangle 6$ . z is unique such that f(z) = z.
  - $\langle 2 \rangle 1. \ f(z) = z$

Proof:

$$f(z) = f\left(\lim_{n \to \infty} x_n\right)$$

$$= \lim_{n \to \infty} f(x_n)$$

$$= \lim_{n \to \infty} x_{n+1}$$

- $\langle 2 \rangle 2$ . If f(w) = w then w = z.
  - $\langle 3 \rangle 1. \| w z \| \le \alpha \| w z \|$

Proof: 
$$||w - z|| = ||f(w) - f(z)|| \le \alpha ||w - z||$$

 $\langle 3 \rangle 2$ . ||w - z|| = 0

PROOF: Otherwise ||w - z|| < ||w - z||.

 $\langle 3 \rangle 3. \ w = z$ 

#### Inner Product Spaces 21.6

Definition 21.6.1 (Inner Product Space). An inner product on a complex vector space V is a function  $\langle , \rangle : V^2 \to \mathbb{C}$  such that:

- $\forall x, y \in V.\langle x, y \rangle = \overline{\langle y, x \rangle}$
- $\forall \lambda \in \mathbb{C}. \forall x, y \in V. \langle \lambda x, y \rangle = \lambda \langle x, y \rangle$
- $\forall x, y, z \in V. \langle x + y, z \rangle = \langle x, z \rangle + \langle y, z \rangle$
- $\forall x \in V. \langle x, x \rangle = 0 \Rightarrow x = 0$

An inner product space or pre-Hilbert space is a complex vector space with an inner product.

**Proposition 21.6.2.**  $\mathbb{C}^n$  is an inner product space under

$$\langle (x_1,\ldots,x_n),(y_1,\ldots,y_n)\rangle = x_1\overline{y_1}+\cdots+x_n\overline{y_n}$$
.

**Proposition 21.6.3.**  $l^2$  is an inner product under

$$\langle (x_n), (y_n) \rangle = \sum_{n=0}^{\infty} x_n \overline{y_n} .$$

 $\langle 1 \rangle 1$ . For all  $(x_n), (y_n) \in l^2$  we have  $\sum_n x_n \overline{y_n} < \infty$   $\langle 2 \rangle 1$ . Let:  $(x_n), (y_n) \in l^2$ 

$$\langle 2 \rangle 1$$
. Let:  $(x_n), (y_n) \in l^2$   
 $\langle 2 \rangle 2$ .  $\sum_{n=1}^N |x_n \overline{y_n}| \leq \left(\sum_{n=1}^\infty |x_n|^2\right)^{1/2} \left(\sum_{n=1}^\infty |y_n|^2\right)^{1/2}$   
PROOF:

OOF:  

$$\sum_{n=1}^{N} |x_n \overline{y_n}| = \sum_{n=1}^{N} |x_n| |y_n|$$

$$\leq \left(\sum_{n=1}^{N} |x_n|^2\right)^{1/2} \left(\sum_{n=1}^{N} |y_n|^2\right)^{1/2} \qquad \text{(Cauchy-Schwarz)}$$

$$\leq \left(\sum_{n=1}^{\infty} |x_n|^2\right)^{1/2} \left(\sum_{n=1}^{\infty} |y_n|^2\right)^{1/2}$$

- $\langle 2 \rangle 3. \sum_{n=1}^{N} x_n \overline{y_n}$  is absolutely convergent.
- $\langle 1 \rangle 2. \ \forall (x_n), (y_n) \in l^2. \langle (x_n), (y_n) \rangle = \overline{\langle (y_n), (x_n) \rangle}$
- $\langle 1 \rangle 3. \ \forall \lambda \in \mathbb{C}. \forall (x_n), (y_n) \in l^2. \langle \lambda(x_n), (y_n) \rangle = \lambda \langle (x_n), (y_n) \rangle$
- $\langle 1 \rangle 4. \ \forall (x_n), (y_n), (z_n) \in l^2. \langle (x_n) + (y_n), (z_n) \rangle = \langle (x_n), (z_n) \rangle + \langle (y_n), (z_n) \rangle$   $\langle 1 \rangle 5. \ \forall (x_n) \in l^2. \langle (x_n), (x_n) \rangle = 0 \Rightarrow (x_n) = 0$

**Proposition 21.6.4.** The space C[0,1] of all continuous functions  $[0,1] \to \mathbb{C}$ is an inner product space under

$$\langle f, g \rangle = \int_0^1 f(t) \overline{g(t)} dt$$
.

**Proposition 21.6.5.** The space  $\mathbb{C}^{mn}$  of all  $m \times n$  complex matrices is an inner product space under

$$\langle A, B \rangle = \operatorname{tr}(B^*A)$$

where  $B^*$  is the conjugate transpose of B.

**Theorem 21.6.6.** Let V be an inner product space. Let  $x, y, z \in V$  and  $\lambda \in \mathbb{C}$ . Then:

- 1.  $\langle x, y + z \rangle = \langle x, y \rangle + \langle x, z \rangle$
- 2.  $\langle x, \lambda y \rangle = \overline{\lambda} \langle x, y \rangle$
- 3.  $\langle x, 0 \rangle = \langle 0, x \rangle = 0$
- 4. If  $\forall w \in V : \langle x, w \rangle = \langle y, w \rangle$  then x = y
- 5.  $\langle x, x \rangle$  is a non-negative real.

PROOF: For part 4, take w = x - y.  $\square$ 

**Theorem 21.6.7** (Cauchy-Schwarz Inequality). Let V be an inner product space. Let  $x, y \in V$ . Then

$$|\langle x, y \rangle| \le \langle x, x \rangle^{1/2} \langle y, y \rangle^{1/2}$$
.

Equality holds if and only if x and y are linearly dependent.

- $\langle 1 \rangle 1$ . If x and y are linearly dependent then  $|\langle x,y \rangle| = \langle x,x \rangle^{1/2} \langle y,y \rangle^{1/2}$ . PROOF: If  $y = \lambda x$  then both sides are equal to  $|\lambda|\langle x, x \rangle$ .
- $\langle 1 \rangle 2$ . If x and y are linearly independent then  $|\langle x, y \rangle| < \langle x, x \rangle^{1/2} \langle y, y \rangle^{1/2}$ .
  - $\langle 2 \rangle 1$ . For any  $\lambda \in \mathbb{C}$  with  $x + \lambda y \neq 0$  we have

$$\langle x, x \rangle + 2\Re(\overline{\lambda}\langle x, y \rangle) + |\lambda|^2 \langle y, y \rangle > 0$$
.

Proof:

$$\begin{split} 0 < \langle x + \lambda y, x + \lambda y \rangle \\ &= \langle x, x \rangle + \overline{\lambda} \langle x, y \rangle + \lambda \langle y, x \rangle + |\lambda|^2 \langle y, y \rangle \\ &= \langle x, x \rangle + 2 \Re(\overline{\lambda} \langle x, y \rangle) + |\lambda|^2 \langle y, y \rangle \\ \langle 2 \rangle 2. \ \text{Let:} \ u = |\langle x, y \rangle| / \langle x, y \rangle \ \text{or} \ u = 1 \ \text{if} \ \langle x, y \rangle = 0 \end{split}$$

- $\langle 2 \rangle 3$ . For any  $t \in \mathbb{R}$ ,

$$\langle x, x \rangle + 2|\langle x, y \rangle|t + \langle y, y \rangle t^2 > 0$$

PROOF: Take  $\lambda = tu$  in  $\langle 2 \rangle 1$ 

 $\langle 2 \rangle 4$ .

$$4|\langle x,y\rangle|^2 - 4\langle x,x\rangle\langle y,y\rangle < 0$$

PROOF: The quadratic  $\langle 2 \rangle 3$  must have negative discriminant.  $\langle 2 \rangle 5.$ 

$$|\langle x, y \rangle| < \langle x, x \rangle^{1/2} \langle y, y \rangle^{1/2}$$

**Theorem 21.6.8.** Every inner product space is a normed space under  $||x|| = \langle x, x \rangle^{1/2}$ .

Proof:

$$\langle 1 \rangle 1$$
. If  $||x|| = 0$  then  $x = 0$ 

$$\langle 1 \rangle 2$$
.  $||\lambda x|| = |\lambda| ||x||$ 

$$\langle 1 \rangle 3. \ \|x + y\| \le \|x\| + \|y\|$$

Proof:

$$||x + y||^2 = \langle x + y, x + y \rangle$$

$$= \langle x, x \rangle + \langle x, y \rangle + \langle y, x \rangle + \langle y, y \rangle$$

$$= ||x||^2 + 2\Re\langle x, y \rangle + ||y||^2$$

$$\leq ||x||^2 + 2|\langle x, y \rangle| + ||y||^2$$

$$< ||x||^2 + 2||x|| ||y|| + ||y||^2$$
 (Cauchy-Schwarz)
$$= (||x|| + ||y||)^2$$

**Theorem 21.6.9** (Parallelogram Law). Let V be an inner product space. Let  $x, y \in V$ . Then

$$||x + y||^2 + ||x - y||^2 = 2||x||^2 + 2||y||^2$$
.

Proof:

$$\begin{aligned} \|x + y\|^2 &= \|x\|^2 + \langle x, y \rangle + \langle y, x \rangle + \|y\|^2 \\ \|x - y\|^2 &= \|x\|^2 - \langle x, y \rangle - \langle y, x \rangle + \|y\|^2 \end{aligned}$$

**Theorem 21.6.10** (Polarization Identity). Let V be an inner product space. Let  $x, y \in V$ . Then

$$4\langle x,y\rangle = \|x+y\|^2 - \|x-y\|^2 + i\|x+iy\|^2 - i\|x-iy\|^2 \ .$$

Proof: Straightforward calculation.