# Mathematics

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## Chapter 1

## **Sets and Functions**

#### 1.1 Primitive Terms

Let there be sets.

Given sets A and B, let there be functions from A to B. We write  $f: A \to B$  iff f is a function from A to B, and call A the domain of f and B the codomain.

Given functions  $f:A\to B$  and  $g:B\to C$ , let there be a function  $gf=g\circ f:A\to C$ , the *composite* of f and g.

#### 1.2 Axioms

#### 1.2.1 Associativity

**Axiom 1.1** (Associativity). For any functions  $f:A\to B,\ g:B\to C$  and  $h:C\to D$  we have

$$h(gf) = (hg)f$$
.

Thanks to this axiom, we shall often omit parentheses when writing the composite of a sequence of functions.

#### 1.2.2 Identity Functions

**Definition 1.2** (Identity Function). For any set A, an *identity function* on A is a function  $i: A \to A$  such that:

- for every set B and function  $f: A \to B$  we have fi = f;
- for every set B and function  $f: B \to A$  we have if = f.

Axiom 1.3 (Identity Functions). Every set has an identity function.

**Proposition 1.4.** Every set has a unique identity function.

Proof:

- $\langle 1 \rangle 1$ . Let: A be a set.
- $\langle 1 \rangle 2$ . A has an identity function.

PROOF: Axiom of Identity Functions

- $\langle 1 \rangle 3$ . For any identity functions i and j on A we have i = j.
  - $\langle 2 \rangle 1$ . Let: i and j be identity functions on A.
  - $\langle 2 \rangle 2$ . i = j

Proof: i = ij = j

**Definition 1.5** (Identity Function). For any set A, let  $id_A$  be the identity function on A.

**Definition 1.6** (Retraction, Section). Let  $r: A \to B$  and  $s: B \to A$ . Then r is a retraction of s, and s is a section of r, iff  $rs = \mathrm{id}_B$ .

**Proposition 1.7.** Let  $f: A \to B$  and  $g, h: B \to A$ . If g is a retraction of f and h is a section of f then g = h.

Proof:

$$g = gid_B$$

$$= gfh$$

$$= id_A h$$

$$= h$$

**Definition 1.8** (Bijection). Let  $f: A \to B$  be a function. We say f is a bijection, and write  $f: A \approx B$ , iff there exists a function  $f^{-1}: B \to A$ , an inverse to f, such that  $f^{-1}f = \mathrm{id}_A$  and  $ff^{-1} = \mathrm{id}_B$ .

Sets A and B are equinumerous,  $A \approx B$ , iff there exists a bijection between them.

Proposition 1.9. The inverse to a bijection is unique.

Proof: From Proposition 1.7.  $\sqcup$ 

#### 1.2.3 The Terminal Set

**Definition 1.10** (Terminal Set). A set T is *terminal* iff, for every set X, there exists exactly one function  $X \to T$ .

**Axiom 1.11** (Terminal Set). There exists a terminal set.

**Proposition 1.12.** If T and T' are terminal sets then there exists a unique bijection  $T \approx T'$ .

Proof:

- $\langle 1 \rangle 1$ . Let: i be the unique function  $T \to T'$
- $\langle 1 \rangle 2$ . Let:  $i^{-1}$  be the unique function  $T' \to T$
- $\langle 1 \rangle 3$ .  $ii^{-1} = id_{T'}$

PROOF: Since there is only one function  $T' \to T'$ .

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 $\langle 1 \rangle 4. \ i^{-1}i = id_T$ 

PROOF: Since there is only one function  $T \to T$ .

**Definition 1.13.** Let 1 be the terminal set. For any set A, let  $!_A$  be the function  $A \to 1$ .

**Definition 1.14** (Element). For any set A, an *element* of A is a function  $1 \to a$ . We write  $a \in A$  for  $a : 1 \to A$ .

Given  $f: A \to B$  and  $a \in A$ , we write f(a) for  $fa: 1 \to B$ .

**Axiom 1.15** (Extensionality). Let A and B be sets and  $f, g : A \to B$ . If  $\forall a \in A. f(a) = g(a)$  then f = g.

#### 1.2.4 The Empty Set

**Axiom 1.16** (Empty Set). There exists a set that has no elements.

#### 1.2.5 Products

**Definition 1.17** (Product). Let A and B be sets. A product of A and B consists of a set  $A \times B$  and functions  $\pi_1 : A \times B \to A$ ,  $\pi_2 : A \times B \to B$ , the projections, such that, for any set X and functions  $f : X \to A$ ,  $g : X \to B$ , there exists a unique function  $\langle f, g \rangle : X \to A \times B$  such that

$$\pi_1 \circ \langle f, g \rangle = f, \qquad \pi_2 \circ \langle f, g \rangle = g.$$

Axiom 1.18 (Products). Any two sets have a product.

**Proposition 1.19.** If P and Q are products of A and B with projections  $p_1: P \to A$ ,  $p_2: P \to B$ ,  $q_1: Q \to A$  and  $q_2: Q \to B$ , then there exists a unique isomorphism  $i: P \approx Q$  such that  $q_1i = p_1$  and  $q_2i = p_2$ .

Proof.

 $\langle 1 \rangle 1$ . Let:  $i: P \to Q$  be the unique function such that  $p_1 i = q_1$  and  $p_2 i = q_2$ .

 $\langle 1 \rangle 2$ . Let:  $i^{-1}: Q \to P$  be the unique function such that  $q_1 i^{-1} = p_1$  and  $q_2 i^{-1} = p_2$ .

 $\langle 1 \rangle 3. \ i^{-1}i = \mathrm{id}_P$ 

PROOF: Each is the unique x such that  $p_1x = p_1$  and  $p_2x = p_2$ .

 $\langle 1 \rangle 4$ .  $ii^{-1} = id_Q$ 

PROOF: Each is the unique x such that  $q_1x = q_1$  and  $q_2x = q_2$ .

**Definition 1.20.** For any sets A and B, we write  $A \times B$  for the product of A and B, and  $\pi_1: A \times B \to A$ ,  $\pi_2: A \times B \to B$  for the projections. Given  $f: X \to A$  and  $g: X \to B$ , we write  $\langle f, g \rangle$  for the unique function  $X \to A \times B$  such that

$$\pi_1 \langle f, g \rangle = f, \qquad \pi_2 \langle f, g \rangle = g.$$

**Definition 1.21.** Given functions  $f:A\to B$  and  $g:C\to D$ , let  $f\times g=\langle f\circ\pi_1,g\circ\pi_2\rangle:A\times C\to B\times D$ .

#### 1.2.6 Function Sets

**Definition 1.22** (Function Set). Let A and B be sets. A function set from A to B consists of a set  $B^A$  and function  $\epsilon: B^A \times A \to B$ , the evaluation map, such that, for any set I and function  $q: I \times A \to B$ , there exists a unique function  $\lambda q: I \to B^A$  such that  $\epsilon \circ (\lambda q \times \mathrm{id}_A) = q$ .

Axiom 1.23 (Function Sets). Any two sets have a function set.

**Proposition 1.24.** If F and G are function sets of A and B with evaluation maps  $e: F \times A \to B$  and  $e': G \times A \to B$ , then there exists a unique isomorphism  $i: F \cong G$  such that  $e'(i \times id_A) = e$ .

#### Proof:

 $\langle 1 \rangle 1$ . Let:  $i: F \to G$  be the unique function such that  $e'(i \times id_A) = e$ .

 $\langle 1 \rangle 2$ . Let:  $i^{-1}: G \to F$  be the unique function such that  $e(i^{-1} \times id_A) = e'$ 

 $\langle 1 \rangle 3$ .  $ii^{-1} = id_G$ 

PROOF: Each is the unique x such that  $e'(x \times id_A) = e'$ .

 $\langle 1 \rangle 4$ .  $i^{-1}i = \mathrm{id}_F$ 

PROOF: Each is the unique x such that  $e(x \times id_B) = e$ .

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#### 1.2.7 Inverse Images

**Definition 1.25** (Pullback). Let  $p:A\to B,\ q:A\to C,\ f:B\to D$  and  $g:C\to D$ . Then we say that  $A,\ p$  and q form the *pullback* of f and g if and only if:

- fp = gq
- For any set X and functions  $x: X \to B$ ,  $y: X \to C$  such that fx = gy, there exists a unique function  $(x,y): X \to A$  such that p(x,y) = x and q(x,y) = y.

We also say p is the pullback of g along f, or g is the pullback of f along g.

$$A \xrightarrow{p} B$$

$$\downarrow f$$

$$C \xrightarrow{g} D$$

**Axiom 1.26** (Inverse Images). Given any function  $f: X \to Y$  and element  $y \in Y$ , then there exists a pullback of f and y.

#### 1.2.8 The Subset Classifier

**Definition 1.27** (Injective). A function  $f: A \to B$  is *injective* iff, for every set X and functions  $x, y: X \to A$ , if fx = fy then x = y.

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**Definition 1.28** (Subset Classifier). A subset classifier consists of a set 2 and an element  $T \in 2$  such that, for any sets A and X and injective function  $j: A \rightarrowtail X$ , there exists a unique function  $\chi: X \to 2$ , the classifying function of j, such that j and  $!_A: A \to 1$  form the pullback of T and  $\chi$ .

$$\begin{array}{ccc}
A & \xrightarrow{!_A} & 1 \\
\downarrow \downarrow & & \downarrow \uparrow \\
X & \xrightarrow{\chi} & 2
\end{array}$$

Axiom 1.29 (Subset Classifier). There exists a subset classifier.

**Proposition 1.30.** If  $T \in 2$  and  $T' \in 2'$  are subset classifiers, then there exists a unique isomorphism  $i : 2 \approx 2'$  such that i(T) = T'.

#### Proof:

⟨1⟩1. Let:  $i:2\to 2'$  be the unique function such that  $\top$  and id<sub>1</sub> form the pullback of  $\top'$  and i

 $\langle 1 \rangle 2$ . Let:  $i^{-1}: 2' \to 2$  be the unique function such that  $\top'$  and  $\mathrm{id}_1$  form the pullback of  $\top$  and  $i^{-1}$ 

 $\langle 1 \rangle 3$ .  $ii^{-1} = id_{2'}$ 

PROOF: Each is the unique x such that  $\top'$  and  $\mathrm{id}_1$  form the pullback of  $\top'$  and x.

 $\langle 1 \rangle 4$ .  $i^{-1}i = id_2$ 

PROOF: Each is the unique x such that  $\top$  and  $\mathrm{id}_1$  form the pullback of  $\top$  and x.

**Definition 1.31.** Let 2 and  $T \in 2$  be the subset classifier.

#### 1.2.9 The Natural Numbers

**Definition 1.32** (Natural Numbers Set). A natural numbers set consists of a set  $\mathbb{N}$ , an element  $0 \in \mathbb{N}$  and a function  $s : \mathbb{N} \to \mathbb{N}$  such that, for any set A, element  $a \in A$  and function  $f : A \to A$ , there exists a unique function  $r : \mathbb{N} \to A$  such that r(0) = a and  $f \circ r = r \circ s$ .

**Axiom 1.33** (Infinity). There exists a natural numbers set.

**Proposition 1.34.** If N,  $0 \in N$ ,  $s: N \to N$  and N',  $0' \in N'$ ,  $s': N' \to N'$  are two natural numbers sets, then there exists a unique isomorphism  $i: N \approx N'$  such that i(0) = 0' and s'i = is.

#### Proof:

 $\langle 1 \rangle 1$ . Let:  $i: N \to N'$  be the unique function such that i(0) = 0' and s'i = is.  $\langle 1 \rangle 2$ . Let:  $i^{-1}: N' \to N$  be the unique function such that  $i^{-1}(0') = 0$  and  $si^{-1} = i^{-1}s'$ .

 $\langle 1 \rangle 3.$   $ii^{-1} = \mathrm{id}_{N'}$ PROOF: Each is the unique x such that x(0') = 0' and s'x = xs'.  $\langle 1 \rangle 4.$   $i^{-1}i = \mathrm{id}_N$ PROOF: Each is the unique x such that x(0) = 0 and sx = xs.

#### 1.2.10 The Axiom of Choice

**Definition 1.35** (Surjective). A function  $f: A \to B$  is *surjective* iff, for every element  $b \in B$ , there exists  $a \in A$  such that f(a) = b.

Axiom 1.36 (Choice). Every surjective function has a section.

#### 1.3 Sections and Retractions

**Proposition 1.37.** Let  $r: A \to B$ ,  $r': B \to C$ ,  $s: B \to A$  and  $s': C \to B$ . If s is a section of r and s' is a section of r', then ss' is a section of r'r.

PROOF: Since  $r'rss' = r'id_Bs' = r's' = id_C$ .

#### 1.4 Isomorphisms

**Proposition 1.38.** For any set A we have  $id_A : A \approx A$  and  $id_A^{-1} = id_A$ .

PROOF: Immediate from the fact that  $id_A id_A = id_A$ .  $\square$ 

**Proposition 1.39.** If  $f : A \approx B$  then  $f^{-1} : B \approx A$  and  $(f^{-1})^{-1} = f$ .

PROOF: Since  $ff^{-1} = id_B$  and  $f^{-1}f = id_A$ .  $\square$ 

**Proposition 1.40.** If  $f: A \approx B$  and  $g: B \approx C$  then  $gf: A \approx C$  and  $(gf)^{-1} = f^{-1}g^{-1}$ .

Proof: From Proposition 1.37.  $\square$ 

#### 1.5 Subsets

**Definition 1.41** (Subset). Let  $i: U \to A$ . Then we say that (U, i) is a *subset* of A iff i is injective.

**Definition 1.42.** Let (U,i) and (V,j) be subsets of A. Then we say (U,i) and (V,j) are equal, and write (U,i)=(V,j), iff there exists an isomorphism  $\phi:U\cong V$  such that  $j\phi=i$ .

#### 1.6 Intersections

**Definition 1.43** (Intersection). Let (U, i) and (V, j) be subsets of a set A. Let  $p: W \to U$  and  $q: W \to V$  form the pullback of i under j. Then the *intersection* of (U, i) and (V, j) is defined to be (W, ip) = (W, jq).

1.7. PULLBACKS

#### 1.7 Pullbacks

#### 1.8 Functions

```
Proposition 1.44. Let f: A \to B. Then f is injective if and only if, for all x, y \in A, if f(x) = f(y) then x = y.
```

```
Proof:
\langle 1 \rangle 1. If f is injective then, for all x, y \in A, if f(x) = f(y) then x = y.
   Proof: Immediate from the definition of injective.
\langle 1 \rangle 2. If \forall x, y \in A. f(x) = f(y) \Rightarrow x = y then f is injective.
   \langle 2 \rangle 1. Assume: \forall x, y \in A. f(x) = f(y) \Rightarrow x = y
   \langle 2 \rangle 2. Let: X be a set and s, t: X \to A
   \langle 2 \rangle 3. Assume: fs = ft
   \langle 2 \rangle 4. \ \forall x \in X. s(x) = t(x)
      \langle 3 \rangle 1. Let: x \in X
      \langle 3 \rangle 2. f(s(x)) = f(t(x))
          Proof: \langle 2 \rangle 3
      \langle 3 \rangle 3. \ s(x) = t(x)
          Proof: \langle 2 \rangle 1
   \langle 2 \rangle 5. s=t
      PROOF: Axiom of Extensionality
```

### 1.9 The Internal Logic

**Proposition 1.45.** Let  $i: U \rightarrow A$  be injective. Let  $\chi: A \rightarrow 2$  be its characteristic function. Then, for all  $a \in A$ , we have  $\chi(a) = T$  if and only if there exists  $u \in U$  such that i(u) = a.

#### Proof:

```
 \begin{array}{l} \langle 1 \rangle 1. \text{ If } \chi(a) = \top \text{ then there exists } u \in U \text{ such that } i(u) = a. \\ \text{PROOF: If } \chi \circ a = \top = \top \circ !_1 \text{ then there exists a unique } u: 1 \to U \text{ such that } i \circ u = a \text{ and } !_U \circ u = !_1. \\ \langle 1 \rangle 2. \text{ For all } u \in U \text{ we have } \chi(i(u)) = \top. \\ \text{PROOF: Since } \chi \circ i = \top \circ !_U. \\ \square \end{array}
```

**Proposition 1.46.** Subsets of a set A are equal if and only if they have the same characteristic function.

PROOF: Follows from the fact that pullbacks are unique up to isomorphism and the uniqueness of the characteristic function.  $\Box$ 

**Proposition 1.47.** There are exactly two subsets of 1.

Proof:

```
\langle 1 \rangle 1. PICK a set E with no elements.
\langle 1 \rangle 2. !_E : E \to 1 is injective.
   PROOF: Vacuously, \forall x, y \in E.!_E(x) = !_E(y) \Rightarrow x = y.
\langle 1 \rangle 3. \ (E,!_E) \neq (1, \mathrm{id}_1)
   PROOF: Since there cannot be an isomorphism 1 \cong E.
\langle 1 \rangle 4. For any subsets (U,i) and (V,j) of 1, if (U,i) \neq (U,i) \cap (V,j) then (U,i) =
   \langle 2 \rangle 1. Let: (U, i) and (V, j) be subsets of 1.
   \langle 2 \rangle 2. Let: p: W \to U and q: W \to V form the intersection of (U,i) and
                    (V,j)
   \langle 2 \rangle 3. Assume: (U, i) \neq (W, k)
   \langle 2 \rangle 4. Let: (U, \mathrm{id}_U) \neq (W, p) as subsets of U.
   \langle 2 \rangle5. Let: \chi_U, \chi_W : U \to 2 be the characteristic functions of (U, \mathrm{id}_U) and
                    (W, p) respectively.
   \langle 2 \rangle 6. \ \chi_U \neq \chi_W
   \langle 2 \rangle 7. Pick x \in U
      PROOF: By the Axiom of Extensionality, there exists x \in U such that
      \chi_U(x) \neq \chi_W(x).
   \langle 2 \rangle 8. \ ix = id_1
   \langle 2 \rangle 9. \ x:1 \cong U
   \langle 2 \rangle 10. \ (U,i) = (1, id_1)
\langle 1 \rangle 5. For any subset (U,i) of 1, either (U,i)=(E,!_E) or (U,i)=(1,\mathrm{id}_1).
   \langle 2 \rangle 1. Let: (U, i) be a subset of 1.
   \langle 2 \rangle 2. Assume: (U, i) \neq (E, !_E)
   \langle 2 \rangle 3. \ (U,i) \neq (U,i) \cap (E,!_E) \text{ or } (E,!_E) \neq (U,i) \cap (E,!_E)
   \langle 2 \rangle 4. (U, i) = (1, id_1) or (E, !_E) = (1, id_1)
      Proof: \langle 1 \rangle 4
   \langle 2 \rangle 5. \ (U,i) = (1, id_1)
      Proof: \langle 1 \rangle 3
```

Corollary 1.47.1. There are exactly two elements of 2.

**Definition 1.48** (Falsehood). Let *falsehood*  $\bot$  be the element of 2 that is not  $\top$ .

**Corollary 1.48.1.** 2 is the coproduct of 1 and 1 with injections  $\top$  and  $\bot$ .

#### 1.10 Functions

**Proposition 1.49.** Let  $f: A \to B$ ,  $g: B \to C$  and  $a \in A$ . Then

$$(g \circ f)(a) = g(f(a))$$
.

PROOF: Immediate from the Axiom of Associativity.

**Proposition 1.50.** For any set A, any function  $1 \to A$  is injective.

1. f is surjective.

PROOF: Since there is only one function  $X \to 1$  for any set X.  $\Box$ 

**Proposition 1.51.** Let  $f: A \to B$ . Then the following are equivalent:

```
2. f is a retraction (i.e. f has a section).
    3. For any set X and functions x, y : B \to X, if xf = yf then x = y.
Proof:
\langle 1 \rangle 1. \ 1 \Rightarrow 2
   PROOF: Immediate from the Axiom of Choice.
\langle 1 \rangle 2. 2 \Rightarrow 3
   \langle 2 \rangle 1. Let: s: B \to A be a section of f.
   \langle 2 \rangle 2. Let: X be a set and x, y : B \to X satisfy xf = yf.
   \langle 2 \rangle 3. \ x = y
       PROOF: x = xfs = yfs = y
\langle 1 \rangle 3. \ 3 \Rightarrow 1
   \langle 2 \rangle 1. Assume: 3
   \langle 2 \rangle 2. Let: b \in B
   \langle 2 \rangle 3. Assume: for a contradiction \forall a \in A. f(a) \neq b
   \langle 2 \rangle 4. Let: \psi_1 : B \to 2 be the characteristic function of b.
   \langle 2 \rangle 5. Let: \psi_2 = \bot \circ !_B : B \to 2
   \langle 2 \rangle 6. \ \forall x \in A. \psi_1(f(x)) = \psi_2(f(x))
       \langle 3 \rangle 1. Let: x \in A
       \langle 3 \rangle 2. \ \psi_1(f(x)) \neq \top
          PROOF: Proposition 1.45, \langle 2 \rangle 3, \langle 2 \rangle 4.
       \langle 3 \rangle 3. \ \psi_1(f(x)) = \bot
       \langle 3 \rangle 4. \ \psi_1(f(x)) = \psi_2(f(x))
   \langle 2 \rangle 7. \ \psi_1 \circ f = \psi_2 \circ f
       PROOF: Axiom of Extensionality
    \langle 2 \rangle 8. \ \psi_1 = \psi_2
```

Corollary 1.51.1. A function is bijective iff it is injective and surjective.

### 1.11 Equalizers

PROOF:  $\langle 2 \rangle 1$  $\langle 2 \rangle 9$ .  $\psi_1(b) \neq \psi_2(b)$ 

 $\langle 2 \rangle 10$ . Q.E.D.

**Theorem 1.52.** Any two functions  $f, g: A \to B$  have an equalizer.

PROOF: Since  $\psi_1(b) = \top$  and  $\psi_2(b) = \bot$ .

Proof: This is a contradiction

```
PROOF: Take the inverse image of \delta_B = \langle \mathrm{id}_B, \mathrm{id}_B \rangle : B \rightarrowtail B^2 and \langle f, g \rangle : A \to B^2. \square
```

#### 1.12 The Empty Set

**Theorem 1.53.** If E is a set with no elements, then E has no proper subsets.

PROOF: A proper subset of E would give a proper subset of 1 that is different from  $(E,!_E)$ .  $\square$ 

**Theorem 1.54.** If E is a set with no elements, then for any set X there exists exactly one function  $E \to X$ .

```
Proof:
\langle 1 \rangle 1. Let: E be a set with no elements.
\langle 1 \rangle 2. Let: X be a set.
\langle 1 \rangle 3. There exists a function E \to X.
   \langle 2 \rangle 1. Let: t: 1 \to 2^X be the name of the characteristic function of id<sub>X</sub>:
                    X \to X.
   \langle 2 \rangle 2. Let: \sigma: X \to 2^X be the lambda of the characteristic function of
                    \delta = \langle \mathrm{id}_X, \mathrm{id}_X \rangle : X \to X \times X.
   \langle 2 \rangle 3. Let: p: P \to E and q: P \to X be the pullback of t \circ !_E and \sigma.
      PROOF: t \circ !_E is vacuously injective.
   \langle 2 \rangle 4. p is injective.
      PROOF: It is the pullback of the injective function \sigma.
   \langle 2 \rangle 5. p is bijective.
   \langle 2 \rangle 6. \ q \circ p^{-1} : E \to X
\langle 1 \rangle 4. For any functions f, g : E \to X we have f = g.
   \langle 2 \rangle 1. Let: f, g : E \to X
   \langle 2 \rangle 2. Let: m: M \to E be the pullback of f and g.
   \langle 2 \rangle 3. (M,m) = (E, \mathrm{id}_E)
      Proof: Since E has no proper subsets.
   \langle 2 \rangle 4. \ m: M \cong E
   \langle 2 \rangle 5. f = g
```

**Corollary 1.54.1.** If E and E' are sets with no elements then there exists a unique isomorphism  $E \cong E'$ .

**Definition 1.55** (Empty Set). Let the *empty set*  $\varnothing$  be the set with no elements.

**Theorem 1.56.** For any set A, if there exists a function  $A \to \emptyset$  then  $A \cong \emptyset$ .

PROOF: If  $f:A\to\varnothing$  then A has no elements, because for any  $a\in A$  we have  $f(a)\in\varnothing$ .  $\square$ 

### 1.13 Universal Quantification

**Definition 1.57.** For any set A, let  $t_A: 1 \to 2^A$  be the name of the characteristic function of  $T \circ !_A: A \to 2$ . Define universal quantification  $\forall_A: 2^A \to 2$  to be the characteristic function of  $t_A$ .

#### 1.14Intersection

**Theorem 1.58.** Let X be a set. There exists a function  $\bigcap : 2^{2^X} \to 2^X$  such that, for all  $S \in 2^{2^X}$  and  $a \in X$ , we have

$$\epsilon(\bigcap S,a) = \top \Leftrightarrow \forall A \in 2^X. (\epsilon(S,A) = \top \Rightarrow \epsilon(A,a) = \top)$$

Proof:

 $\langle 1 \rangle 1$ . Let: X be a set.

 $\langle 1 \rangle$ 2. Let:  $\phi_2: X \to 2^{2^X}$  be the lambda of  $\epsilon: 2^X \times X \to 2$ 

 $\langle 1 \rangle 3$ . Let: F be the function

$$2^{2^X} \times X \xrightarrow{\langle \operatorname{id}_{2^{2^X}}, \phi_2 \rangle} 2^{2^X} \times 2^{2^X} \xrightarrow{\cong} (2 \times 2)^{2^X} \xrightarrow{\Rightarrow} 2^{2^X} \xrightarrow{\forall} 2$$

 $\langle 1 \rangle 4$ . Let:  $\bigcap$  be the lambda

#### 1.15Union

**Theorem 1.59.** Any two subsets of a set have a union.

Proof:

 $\langle 1 \rangle 1$ . Let: A and B be subsets of X

 $\langle 1 \rangle$ 2. Let:  $\chi_A \in 2^X$  be the name of the characteristic function of A.

 $\langle 1 \rangle 3$ . Let:  $t_X \in 2^X$  be the name of  $\top \circ !_X : X \to 2$ 

 $\langle 1 \rangle$ 4. Let: C be the pullback of  $t_X$  and  $\chi_A \Rightarrow -: 2^X \to 2^X$ 

 $\langle 1 \rangle$ 5. Let: D be the pullback of  $t_X$  and  $\chi_B \Rightarrow -$ 

 $\langle 1 \rangle 6$ .  $\bigcap (C \cap D)$  is the union of A and B.

**Theorem 1.60.** Any two sets have a coproduct.

Proof:

 $\langle 1 \rangle 1$ . Let: X and Y be sets.

 $\langle 1 \rangle 2$ . Let:  $\sigma_X : X \to 2^X$  be the lambda of the characteristic function of  $\langle \mathrm{id}_X, \mathrm{id}_X \rangle : X \to X \times X$ 

 $\langle 1 \rangle 3$ . Let:  $\chi_0 : 1 \to Y$  be the characteristic function of the unique function

 $\langle 1 \rangle 4$ . Let:  $i_X = \langle \sigma_X, \chi_0 \circ !_X \rangle : X \to 2^X \times 2^Y$  $\langle 1 \rangle 5$ . Let:  $i_Y : Y \to 2^X \times 2^Y$  be defined similarly.

 $\langle 1 \rangle 6$ .  $i_X$  and  $i_Y$  are monic.

 $\langle 1 \rangle 7$ .  $\varnothing$  is the pullback of  $i_X$  and  $i_Y$  (i.e.  $(X, i_X) \cap (Y, i_Y) = \varnothing$ ).  $\langle 1 \rangle 8$ . Let:  $j: Z \to 2^X \times 2^Y$  be the union of  $i_X$  and  $i_Y$ 

 $\langle 1 \rangle 9$ . Z is the coproduct of X and Y.