



7.

There are three foreign keys:-

1. Attribute serial-no of relation OPTION that references relation CAR
2. The Attribute salesperson-id of relation SALE ^{refers} that relation SALESPERSON
3. The attribute serial-no of relation SALE that references relation CAR

SALE				SALESPERSON		
Salesperson-id	Serial-no	Date	Sale-price	Salesperson-id	name	phone
1	99	24-1-20	100	1	RK	9806391416
1	100	25-1-20	90	2	PK	7286009239
2	50	1-1-20	150			

If we try to

INSERT INTO SALE VALUES (3, 44, 1-2-20, 200)

This is violation of referential integrity constraints
Since salesperson-id 3 not exists

INSERT INTO SALE VALUES (2, 66, 2-2-20, 400)

This is not violating because salesperson-id 2 exists

10.

- a) one possible set of operations for the following update is the following:

```
INSERT < FNO, LNO, DT, SEAT_NO, CUST_NAME,
        CUST_PHONE > INTO SEAT-RESERVATION;
```

modify the LEG-INSTANCE TUPLE with the condition:

(FLIGHT-NUMBER = FNO AND LEG-NUMBER = LNO AND
DATE = DT) by setting NUMBER-OF-AVAILABLE-SEATS
= NUMBER-OF-AVAILABLE-SEATS - 1;

These operations should be repeated for each LEG of the flight on which a reservation is made.

This assumes that the reservation has only one seat. More complex operations will be needed for a more realistic reservation that may reserve several seats at once.

- b) we would check that NUMBER-OF-AVAILABLE-SEATS on each LEG-INSTANCE of flight is greater than 1 before doing any reservation, and that the SEAT-NUMBER being reserved in SEAT-RESERVATION is available.

- c) The INSERT operation into SEAT-RESERVATION will check all the key, entity integrity, referential integrity constraints for the relation. The check that NUMBER-OF-AVAILABLE-SEATS on each LEG-INSTANCE of the flight is greater than 1 does not fall into any of the above types of constraints (It is a general semantic integrity constraint).

12. a) Names of parts that cost less than 20.00

SQL > SELECT * FROM PARTS WHERE Price < 20;

Algebraic expression

temp1 $\leftarrow \sigma_{\text{Price} < 20}(\text{PARTS})$

Results $\leftarrow \pi_{\text{Pno}, \text{Pname}, \text{Qoh}, \text{Price}, \text{Qlevel}}(\text{temp1})$

b) Names and cities of employees who have taken orders for parts costing more than 50.00

SELECT Ename, City FROM EMPLOYEES e, ZIP_CODES z, ORDERS o, ODETAIL o1, PARTS p WHERE p.Price > 50 AND p.Pno = o1.Pno AND o.Ono = o1.Ono AND e.Eno = o.Eno AND e.Zip = z.Zip;

g) Names of customers who have placed exactly two orders

SELECT Cname, count(Ono) as total FROM CUSTOMERS c, ORDERS o WHERE c.Cno = o.Ono AND total = 2;

- c) pairs of customer number values of customers who live in the same ZIP code

```
SELECT c.cno, cl.cno FROM customers c, CUSTOMERS cl
WHERE c.zip = cl.zip AND c.cno < cl.cno ;
```

- d) Names of customers who "ordered parts from employees living in wichita

```
SELECT DISTINCT cname FROM CUSTOMERS c, ORDERS o,
EMPLOYEES e, ZIP_CODES z. WHERE c.cno = o.ono AND
c.zip = z.zip AND z.city = 'wichita' ;
```

- e) names of customers who have ordered parts costing less than 20

```
SELECT DISTINCT cname FROM customers c, PARTS p,
ORDERS ODETAILS o, ORDERS ol WHERE WHERE p.Price < 20
AND p.pno = o.pno AND o.ono = ol.ono AND
c.cno = ol.cno ;
```

- (d) f) Names of customers who have not placed any order

```
SELECT cname, count(ono) as total FROM
CUSTOMERS c, ORDERS o. WHERE c.cno = o.ono
AND total = 0 ;
```