6. Unification is a process of making two different logical atomic expressions identical by finding a substitution

Unification depends on the substitution process. It takes two literals as input and makes them identical using substitution.

Let  $\psi_1$  and  $\psi_2$  be two atomic sentences and  $\sigma$  be a unifier such that  $\psi_1\sigma = \psi_2\sigma$ , then it can be expressed as (a)

unity (4, 142) (1) 11.000 319 d

Example: Find the MGW for unify (king (n), king (John) }

Let  $\psi_1 = \ker_q(n)$ ,  $\psi_2 = \ker_q(\operatorname{John})$ 

substitution 0= {(Tohn (n) } is a unitien for these atoms and applying this substitution, and both expressions will be identical.

The UNIFY Algorithm II used for Unification, which takes two atomic sentences and returns a unifier for those sentences (if any exist) unification is a key component of all first order (order) inference algorithms

It returns fail if the expressions do not match with each other.

The substitution variables are called Most

a segue at another required

unification Algorithmi-

Algorithm: unity (41, 42)

step-1: If  $\psi_1$  or  $\psi_2$  is a variable or constant

- a) If 4, or \$1 are identical then return NIL.
- b) fise if  $\psi_1$  is a variable,

a. then if  $\psi_1$  occurs in  $\psi_2$ , then return

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b. εue return { (42/41) }

- c) fixe if Ψ<sub>2</sub> is a variable,
  a. If Ψ<sub>2</sub> occurs in Ψ, then return FAILURE
  b. the return {(Ψ1/Ψ2)}
- d) fise return FAILURE

step-2; If the initial predicate symbol in  $\psi_1$  and  $\psi_2$  are not same theoreturn FAILURE.

1tep-3: It 4, and 42 have a different no of anguments, then return FAILURE

step-4: set substitution set (subst) to NIL.

step T: For i=1 to the no of elements in \$1

- a) call unify function with the 1th element of  $\psi_1$  and  $i^{th}$  element of  $\psi_2$  and put the result into s.
- b) If s=failure then returns FAILURE
- c) It s = NIL then do

a. Apply 1 to the remainder of both Ll and L. b. SUBST = APPEND (S, SUBST)

step-6 - Return SUBST

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Example
-> UNIFY (knows (Richard, xi), knows (Pichard, John))
   Here 41 = knows(Richard, n)
        42 = knows ( Richard, John)
   So =) { knows (Richard, X); knows (Richard, Tohn)}
    JURST 0 = { John (x }
   Si =) { knows (Richard, John); knows (Richard, John)}
   successfully unified
     unifier = { John /x }
       { p(x,x) and p(z, t(z)) }
   Here \psi_1 = \{ p(x_1 x) \}
         \Psi_2 = P(2, f(2))
   So =) {P(x,x), P(2,f(2))}
    SUBJT 0= { x/2}
            SI =) {P(2,2), P(2, f(2))}
    SUST 0 = { P(2)/2}, unification failed.
```