

# Optimization Problem: Mixed BEV and FCEB Fleet for GTFS Routes

## 1. Objective Function

Minimize the total operational cost:

$$C_{\text{total}} = C_{\text{fleet}} + C_{\text{energy}} + C_{\text{infrastructure}} + C_{\text{maintenance}}$$

Where:

$$\begin{aligned} C_{\text{fleet}} &= \sum_{b \in BEVs} C_{\text{BEV.cost}} + \sum_{b \in FCEBs} C_{\text{FCEB.cost}} \\ C_{\text{energy}} &= \sum_{t \in T} \left( \sum_{b \in BEVs} x_{\text{charging},b,t} \cdot C_{\text{electricity}} + \sum_{b \in FCEBs} x_{\text{refueling},b,t} \cdot C_{\text{hydrogen}} \right) \\ C_{\text{infrastructure}} &= \sum_{l \in Chargers} x_{\text{opp.chargers},l} \cdot C_{\text{opp.cost}} + \sum_{l \in RefuelingStations} x_{\text{refueling.stations},l} \cdot C_{\text{refuel.cost}} \\ C_{\text{maintenance}} &= \sum_{b \in Buses} C_{\text{maintenance},b} \end{aligned}$$

## 2. Constraints

### 2.1 Route Coverage

Ensure every route is served by at least one bus:

$$\sum_{b \in BEVs \cup FCEBs} x_{\text{on.route},r,b,t} \geq 1, \quad \forall r \in R, \forall t \in T$$

### 2.2 BEV Energy Balance

Ensure energy balance for BEVs:

$$x_{\text{charging},b,t} + \sum_{r \in R} x_{\text{on.route},r,b,t} \cdot C_{\text{energy.use},r} \leq C_{\text{charge.rate}}, \quad \forall b \in BEVs, \forall t \in T$$

### 2.3 BEV Range Constraint

Ensure BEVs can complete assigned routes based on battery capacity:

$$C_{\text{route.distance},r} \leq x_{\text{battery.size},b}, \quad \forall b \in BEVs, \forall r \in R$$

## 2.4 FCEB Hydrogen Balance

Ensure hydrogen balance for FCEBs:

$$x_{\text{refueling},b,t} + \sum_{r \in R} x_{\text{on.route},r,b,t} \cdot c_{\text{energy.use},r} \leq c_{\text{refuel.rate}}, \quad \forall b \in \text{FCEBs}, \forall t \in T$$

## 2.5 FCEB Range Constraint

Ensure FCEBs can complete assigned routes based on hydrogen tank capacity:

$$c_{\text{route.distance},r} \leq x_{\text{hydrogen.tank.size},b}, \quad \forall b \in \text{FCEBs}, \forall r \in R$$

## 2.6 One Task Per Bus

A bus can perform only one task (route, charging, refueling) at a time:

$$\sum_{r \in R} x_{\text{on.route},r,b,t} + x_{\text{charging},b,t} + x_{\text{refueling},b,t} \leq 1, \quad \forall b \in \text{BEVs} \cup \text{FCEBs}, \forall t \in T$$

## 2.7 Depot Charging Capacity

Ensure depot chargers are not overloaded:

$$\sum_{b \in \text{BEVs}} x_{\text{charging},b,t} \leq x_{\text{depot.chargers}}, \quad \forall t \in T$$

## 2.8 Opportunity Charging Capacity

Ensure opportunity chargers are not overloaded:

$$\sum_{b \in \text{BEVs}} x_{\text{charging},b,t} \leq x_{\text{opp.chargers},l}, \quad \forall l \in \text{Chargers}, \forall t \in T$$

## 2.9 Refueling Station Capacity

Ensure refueling stations are not overloaded:

$$\sum_{b \in \text{FCEBs}} x_{\text{refueling},b,t} \leq c_{\text{refuel.capacity},l}, \quad \forall l \in \text{RefuelingStations}, \forall t \in T$$

## 2.10 Fleet Size

Ensure the total fleet size matches the number of buses available:

$$\sum_{b \in \text{BEVs}} x_{\text{fleet.BEVs}} + \sum_{b \in \text{FCEBs}} x_{\text{fleet.FCEBs}} = \text{Total Fleet Size}$$

## 3 Time Complexity Analysis

The time complexity of the column generation framework is influenced by three primary components: the Restricted Master Problem (RMP), the Pricing Problem, and the iterative nature of the column generation process.

### 3.1 Restricted Master Problem (RMP)

The RMP involves solving a linear programming problem over a subset of variables, referred to as columns. The computational complexity of solving a linear program depends on the number of variables ( $C$ ) and constraints. For the RMP,  $C$  represents the number of active columns, which corresponds to feasible bus-route assignments. Standard solvers, such as those based on the simplex or interior-point methods, exhibit a worst-case complexity of:

$$O(C^3)$$

This cubic growth arises due to the linear programming solver's reliance on matrix operations, such as factorization.

### 3.2 Pricing Problem

The Pricing Problem is responsible for identifying new columns (bus-route assignments) to improve the RMP solution. This involves evaluating all feasible bus-route combinations while adhering to constraints such as timing, energy capacity, and charging/refueling limits. The complexity of this step scales linearly with the number of buses ( $B$ ) and routes ( $R$ ), resulting in a computational complexity of:

$$O(B \cdot R)$$

per iteration.

### 3.3 Column Generation Iterations

Column generation operates iteratively, alternating between solving the RMP and the Pricing Problem until convergence is achieved. The number of iterations required ( $I$ ) typically depends on the problem's structure and the rate of convergence. Empirical evidence suggests that  $I$  often grows logarithmically with the number of routes, i.e.,:

$$O(\log R)$$

This behavior is attributed to the diminishing marginal impact of adding new columns as the solution approaches optimality.

### 3.4 Overall Time Complexity

The overall complexity of the column generation framework can be expressed as:

$$O(I \cdot (C^3 + B \cdot R))$$

where:

- $I$  is the number of iterations.
- $C$  is the number of columns in the RMP, which increases with the number of routes.
- $B$  is the number of buses.
- $R$  is the number of routes.

For large-scale problems,  $C$  is proportional to  $R$ , and  $I$  is approximated as  $O(\log R)$ . Therefore, the time complexity simplifies to:

$$O(\log R \cdot (R^3 + B \cdot R))$$

### 3.5 Dominant Term

Among the components, the  $R^3$  term from solving the RMP typically dominates for large  $R$ , making the overall framework computationally intensive for problems with many routes. While the Pricing Problem contributes linearly to the complexity ( $B \cdot R$ ), its impact is overshadowed by the cubic growth of the RMP.

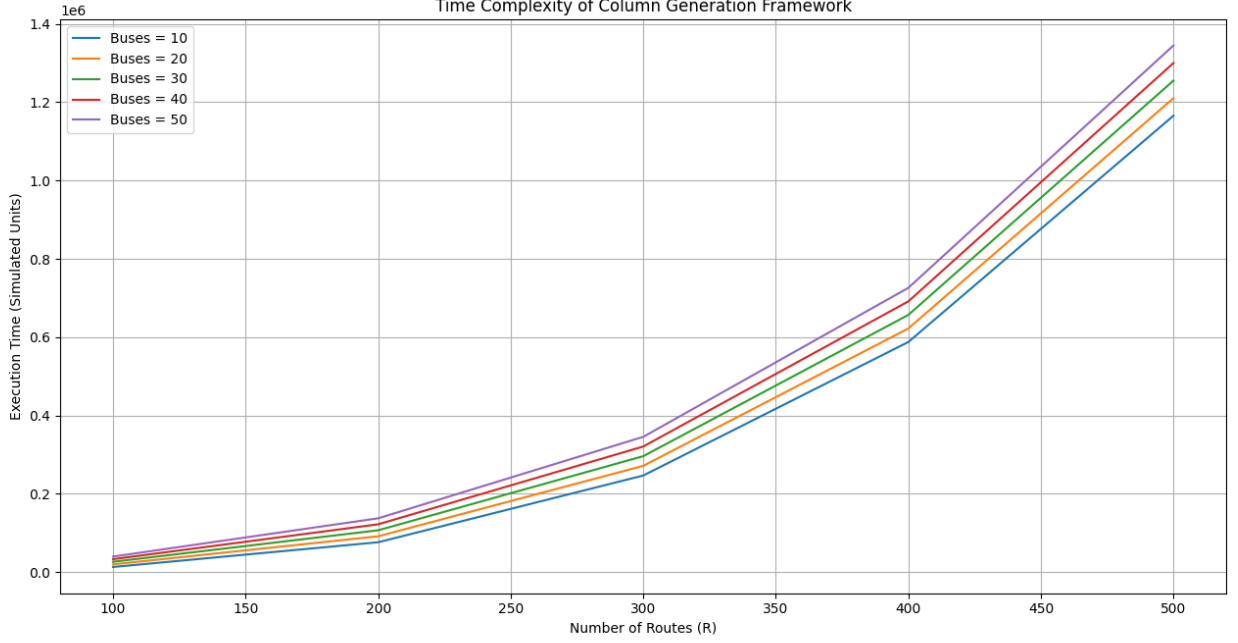


Figure 1: Time Complexity Graph for Column Generation

## Complexity Analysis of the Simulated Annealing Algorithm

- **Iterations:**
  - The algorithm runs for `max_iterations` (let's call it  $I$ ) or until the temperature drops below a certain threshold.
  - **Complexity:**  $O(I)$
- **Neighbor Generation (`generate_neighbor`):**
  - This involves choosing a random route and assigning a new bus.
  - **Complexity:**  $O(1)$ , as it performs a constant-time operation for each route.
- **Cost Calculation (`calculate_cost`):**
  - This is the most computationally intensive function. It involves:
    - \* Iterating over all routes to calculate operational costs and check constraints.

- \* For each bus, summing up energy consumption across routes for energy balancing.
- \* Checking charger and refueling station constraints.
- Let:
  - \*  $R$  = Number of routes.
  - \*  $B$  = Number of buses.
  - \*  $O$  = Number of opportunity chargers.
  - \*  $F$  = Number of refueling stations.
- **Complexity:**
  - \* Route iteration:  $O(R)$ .
  - \* Charger utilization:  $O(B)$ .
  - \* Energy balancing:  $O(B \times R)$  (in the worst case, every bus checks every route).
  - \* Opportunity charger checks:  $O(R \times O)$ .
  - \* Refueling station checks:  $O(R \times F)$ .
- **Combined complexity:**  $O(R \times (B + O + F))$
- **Acceptance Probability Calculation:**
  - Exponential computation and random number generation are  $O(1)$ .