QUIZ 2_ADA 2021 (FS+IoT+GT)

Points: 13/15

In Branch and Bound method (1/1 Point)

- One divides a large problem into few smaller ones, and this is the branching part.
- While the conquering part is done by estimating how good a solution, we can get from each smaller problem.
- Both a and b
- None



2

The order of the worst case time complexity of backtracking algorithms is: (0/1 Point)

- N^3

2^N ✓
N log log N
3
IN TSP branch and Bound, Matrices are reduced by- (1/1 Point)
Row reduction
Column Reduction
■ Both ✓
None
4
Choose the correct statements(Multiple correct) (2/2 Points)
P problems can be solved in polynomial time on deterministic computers
NP problems can be solved in polynomial time on non deterministic computers
NP problems can be solved in polynomial time on deterministic computers
✓ All NP Hard problems lie in set NP. ✓
All NP complete problems lie in NP 🗸

5

○ V×E

The goal of branch and bound is to find a value x thatt real-valued function f(x), called an objective function among some admissible or candidate solutions(multiple correct) (1/1 Point)	
✓ Maximizes ✓	
✓ Minimizes ✓	
Normalizes	
Calculates	
6	
In n-queen problem, the backtrack step is: (1/1 Point)	
to place the queen to next available space	
\bigcirc to go to the previously placed queen and place it to next available space \checkmark	
to go to the first queen and place it to the next available space	
to place the queen to the last available space	
7	
The chromatic number of a complete graph $G=(V, E)$, will be: (1/1 Point)	
∨ ✓	
○ V^2	

2* E

8

Choose the correct statement from the following: (1/1 Point)

Branch and Bound is more efficient than backtracking.
Branch and bound is not suitable where a greedy approach is not applicable.
Branch and Bound divides the problem into at least 2 new restricted sub problems 🗸
Backtracking divides the problem into at least 2 new restricted sub problems.
9
Choose one or more correct statement/s related with Hamiltonian Circuit algorithm: (1/1 Point)
✓ Hamiltonian circuit may or may not exist in a graph ✓
Hamiltonian circuit finding depends upon the choice of initial vertex
Algorithm searches the start/ end vertex irrespective of what we give as initial vertex
In a complete graph, there would not be any backtrack step required to find the circuit

Tractable problems are the problems (2/2 Points)

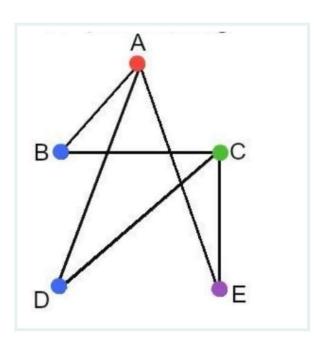
That can be solved in polynomial time \checkmark

That can not be solved in polynomial time

- Problems cannot be solved by any algorithm
- Problems that lie under the category of decision problems \checkmark



11



The chromatic number of the following graph would be: (0/1 Point)

- 2 <
- 3

12

In n-queen problem, the condition to check whether two queens at positions (i, j) and (k, l) are not attacking each other is/are: (1/1 Point)

i ≠ k and j ≠ l

$i \neq k, j \neq l, i \neq j \text{ and } k \neq l$
$i \neq k, j \neq l$ and $i-k \neq j-l $
$i+k\neq j+1 $ \checkmark
13
To calculate Upper bound in knapsack Problem: (1/1 Point)
Fractional weights are considered.
● Fractional weights are not considered. ✓
Can be calculated both ways.

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