

# Design and Analysis of Algorithms

## 6.6 General Functions

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## 7.6.2

- Let  $f(k)$  be the running time of the  $k^{th}$  operation.
- Consider algorithms with the following running times:
  - $f(k)$  = the largest power of 2 that divides  $k$
  - $f(k) = k$  if  $k$  is a power of 2,  $f(k) = 1$
- Calculate amortized running time for each algorithm.
- $f(k)$  = the largest power of 2 that divides  $k$

1, 2, 1, 4, 1, 2, 1, 8, 1, 2

$$\frac{n}{3} + \sum_{k=0}^n \frac{k}{2}$$

...?

$n$  amortized runtime

- $f(k) = k$  if  $k$  is a power of 2,  $f(k) = 1$

1, 2, 3, 4, 5, 6, 7, 8

1, 1, 3, 1, 5, 6, 7, 1

$$\frac{n}{1} + \sum_{k=0; k \neq 2}^n k$$

$n \log(n)$ ?

## Solution

- $f(k)$  = the largest power of 2 that divides  $k$

$$1 \frac{n}{2} + 2 \frac{n}{4} + 4 \frac{n}{8} + 2 \frac{n}{4} + \dots + \frac{2^{\log(n)}}{2} * \frac{n}{2^{\log(n)}} + n$$

$$\frac{n}{2} + \frac{n}{2} + \frac{n}{2} + \dots + \frac{n}{2} + n$$

$$\frac{n}{2} \log(n) + n$$