



School of
Engineering

Bachelor's thesis
Information Science

Design and Implementation of an Alternative to SSH

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Abstract

As [Secure Shell \(SSH\)](#) is an old application that has loads of features. Not all of those features are used in everyday business. In this thesis a simpler protocol prototype that can do the same as [SSH](#) was designed and implemented. The developed solution can replace [SSH](#) in its core feature: Connecting to a [shell](#) on a remote system.

This project developed an application called “oh-my-gosh” that is capable of the same core functionality that [SSH](#) provides. This includes a client application called “gosh” and a server counterpart with the name “goshd”. The latter was designed to be run as a background process on a [Linux](#) system. This solution uses a secure connection as channel to provide more safety in the entire process. A user can authenticate itself on the remote system either via a password or using public key cryptography. All the usual work flows are possible on the [shell](#) such as:

- Navigating through the file system
- Running scripts and applications
- Using applications that use [ncurses\(3X\)](#) (2019)

The solution relies on a number of [Unix](#) specific technologies like [pseudoterminals \(ptys\)](#) and [Pluggable Authentication Module \(PAM\)](#), which is used in [login\(1\)](#) (2012) as well, to realize its use cases.

Preface

The **SSH** protocol ([Moorer 1971](#), [Bider & Baushke 2012](#), [Baushke 2017](#), [Bider 2018a,b](#)) is a system that allows a user to log in on a remote machine and perform tasks on that remote machine via a **Command Line Interface (CLI)**. **SSH** is widely known and used in everyday tasks. However: It is now over twelve years old in its current form. One of the problems with **SSH** is its complexity, both in the initial phase when key material is exchanged, but also later, for example because the server-side has to solve whether to return a character that has been sent to it or not (echo).

The goal of this work is a radically simplified protocol, which in its functions is similar to **SSH** (N.B. the similarity concerns the functions, not necessarily the protocol details). I develop the protocol, as well as a client and a server - all in [the Go/Golang programming language \(Go\)](#). I demonstrate that the software can replace **SSH** by showing that it can handle several common use cases, among them:

- Interactive session
- Rsync with my solution as transport protocol

This bachelor thesis was proposed by Dr. Stephan Neuhaus ([Neuhaus 2018](#)) and aroused my interest as it is a challenge in the domain of information security and will produce a palpable result.

On this note I would like to thank Dr. Stephan Neuhaus for helping me along the way of this Bachelors thesis.

As an additional remark, I would like to point out that the team for this Bachelors thesis originally consisted of two people, but after some months of uncertainty, Mr Schwarz decided to opt out of the project as he said he underestimated the workload from the modules he booked on top of the Bachelors thesis. This happened halfway through the project and had great impact on the project itself: The expected work from him had to be done by me, Raphael Emberger. This also meant less time for designing and implementing all the features described in the tasks (see [1.4](#)). Fortunately Dr. Neuhaus adapted those to fit a one-man project.

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1. Introduction

There was no thesis done on this subject that could have been used as reference. There are however several software projects that deal with a similar problem.

1.1. OpenSSH

The most noteworthy work to mention is of course [SSH](#) itself. [OpenSSH \(1999\)](#) is the name of the open source project which provides millions of administrators and developers with the ability to securely connect to a remote host. It replaces the up until then widely used protocols like [telnet\(1\) \(1994\)](#) (see [1.2](#)) and [rlogin\(1\) \(1999\)](#)/[rsh\(1\) \(1999\)](#) (see [1.3](#)).

[SSH](#) uses an own protocol to secure the communication channel between two peers and has earned itself a spot on the low end of the [port](#) table: It occupies [port 22](#).

[SSH](#)'s features can be used very flexibly: After it builds up a secure connection between a client and a server, it can be used to remotely log in and use a [terminal](#) on that machine. It can also forward traffic on local ports to the remote host through the secure channel. This is also used by third party programs such as [rsync\(1\) \(2018\)](#).

When it comes to the log in procedure itself, [SSH](#) allows for standard user [login](#) using the [Application Programming Interface \(API\)](#) of [PAM](#). Another feature is white listing of clients via their public keys, which stops intrusion attempts via hijacked user-password-credentials.

After a secure connection could be established, there are multiple possibilities to use the opened channel. One is to forward the [Graphical User Interface \(GUI\)](#) of a remote program to the client. Another one is to use this channel to tunnel more connections through it: For example can the traffic of an application which uses a specific [port](#) be forwarded to the remote host. This can obscure and secure this traffic between the host and the server.

1.2. Telnet

Telnet ([C. Stephen 1969](#), [Postel & Reynolds 1983](#)) is an old (1969) and deprecated communication protocol which does not feature any security. However, in other implementations, [Telnet Secure \(TELNETS\)](#) was proposed, which features encryption over the communication channel. Telnet still has 23 as its very own [port](#) assigned to it.

It is also to mention that there is an application called Go-Telnet ([Krempeaux 2016](#)). This is a [TELNETS](#) supporting client-server-application which has been implemented in [Go](#).

1.3. Berkeley r-Commands

The Berkley r-commands are a set of commands to do certain tasks on remote hosts. Those tasks are similar to their counterparts without a leading "r".

- [rlogin\(1\) \(1999\)](#)

This command connects to the host and performs a [login\(1\) \(2012\)](#) command, which includes authentication and if successful, spawning a user [shell](#).

- [rsh\(1\) \(1999\)](#)

`rsh` executes a command on the remote host. If no command is specified, the user gets logged in on the host with [rlogin\(1\) \(1999\)](#).

- *rexec(1)* (1996)

With this command, the user can log in to a remote machine and execute one command.

- *rcp(1)* (1999)

Using this command gives the user the ability to copy from and to a remote host.

- *rwho(1)* (1996)

This command tells the user what users are currently logged in on the remote machine.

- *rstat(1)* (1996)

`rstat` displays file system information from remote hosts.

- *ruptime(1)* (1996)

With this command, the user can see the [uptime](#), number of logged in users and current work load of the remote machine.

It is to be noted that none of those commands support securing the communication channel in any way. This is also the reason why they got replaced by [SSH](#), as the latter does secure it's tunnel between a client and a server.

1.4. Task

This project's objective is to design and implement a prototype for an alternative to [SSH](#). It has to be able to provide the user with the ability to have an interactive session on a remote machine. The solution has to be able to connect to a remote server and use a [shell](#) there. Both client and server-side have to be designed and implemented.

The official formulation of the tasks can be found in the appendix [A.1.1](#).

The chapter "Design" (see [2](#)) is meant to give a detailed overview of how the solution initially was envisioned and designed. In the course of the project's development, many aspects and views changed, as can be read in the minutes (see [A.1.3](#)). To clarify these changes in detail is the "Implementation" chapter's objective. In there, another overview has been worded, which should give sufficient insight on the project's development.

This thesis has been worded with technically literate readers in mind. However: For core concepts and special terms, a glossary can be found at [6](#). Used acronyms are listed in [6.4](#).

2. Design

To create a client-server protocol for remote [login](#) and interactive sessions, a clear cut architecture is mandatory. The flow of a typical use case should look like this:

1. The server listens for incoming connection.
2. The client dials server.
3. The server spawns the users [login shell](#) and forwards all traffic between [shell](#) and Client.
4. The client uses the [shell](#).
5. The client terminates the session.
6. The server listens for new incoming connections.

However, there are multiple security concerns to be satiated:

- The connection between the client and server has to be secured from exposure to or manipulation from third parties.
- The client must authenticate itself for a user of the remote system with the appropriate credentials.
- The server has to drop privilege after a successful [login](#) to prevent privilege escalation.
- The server has to spawn the [login shell](#) of the logged in user with the appropriate rights.

Furthermore, this design does not allow for multiple sessions to be run in parallel. Therefore it was decided to run a new kernel-level thread to handle everything after the connection has been established. This thread could also drop its privileges after the [login](#) succeeded.

Another action the server would have to take is to properly manage the session. That means the server has to make sure the [login](#) time stamp, who logged in, what device was used and other parameters are stored on the system. This is to act according to the [Unix](#) specification. This [login](#) accounting will allow for the logged in user to be displayed with commands like `w`.

This led to the following general flow of actions:

1. The server listens for incoming connections.
2. The client dials server.
3. The server handles the established connection.
4. The server sets up a secure connection (see [2.2](#)).
5. The server requests necessary environment variables from the client.
6. The client responds with the values of said variables.
7. The server initiates a [login](#) procedure.
8. The client authenticates itself (see [2.3](#)).
9. The server gathers essential data about the logged in user (see [2.4](#)).
10. The server performs the necessary post-[login](#) actions that are expected in a [Unix](#) environment (see [2.6](#)).

11. The server drops privilege to that of the logged in user (see 2.7).
12. The server spawns the user's [login shell](#) (see 2.5) with the same credentials and forwards all traffic between [shell](#) and client.
13. The client uses the [shell](#) on the remote machine.
14. The client terminates the session.
15. The server terminates.

The sequence diagram in figure 2.1 is an illustration of how the flow of actions was envisioned.

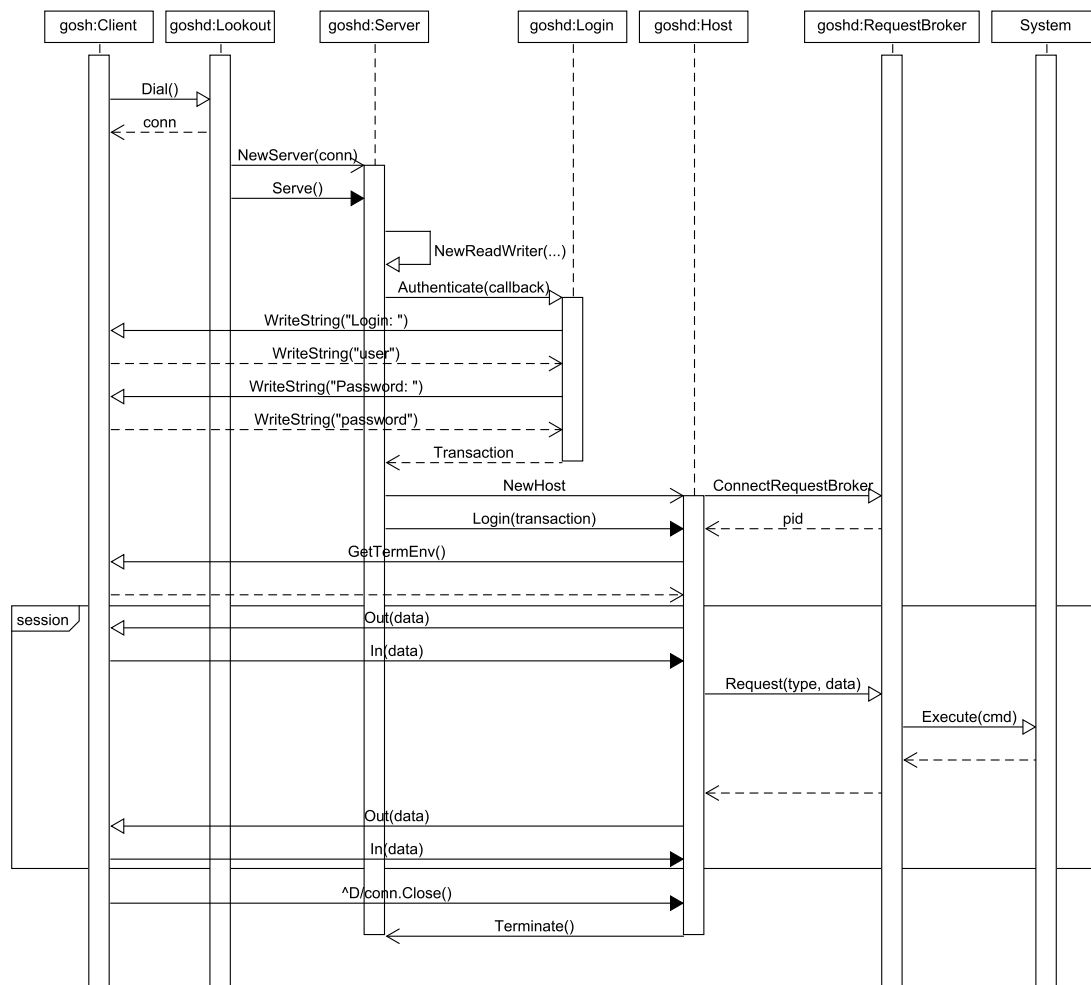


Figure 2.1.: Sequence diagram draft.

2.1. Implementation Language

In the beginning of the project, Dr Neuhaus suggested the use of [Go](#) as a modern low-level language over interpreted languages for considerations of security. He emphasized that the use of other low-level languages (like [the C programming language \(C\)](#) or [the C++ programming language \(C++\)](#)) was permissible. In the end the project was implemented in [Go](#) as suggested.

However, this led to a few problems in the implementation process (See 3.1).

2.2. Secure Connection

When building an application that communicates via the network, certain security measures are mandatory to ensure a secure communication. If such measures are not taken, the communication between the client and the server can be fully read and even altered by a third party. To prevent this, the communication can be encrypted with [Transport Layer Security \(TLS\)](#) (originally known as [Secure Sockets Layer \(SSL\)](#)). Today's state-of-the-art is [TLS 1.3](#) ([Rescorla 2018](#)).

When connecting, the client checks the server's [X.509](#)-certificate to authenticate the peer. Optionally, the client can also authenticate itself to the server. After this, the two start an encrypted channel by for example using the [Diffie-Hellman key exchange](#) or letting the server decrypt a random secret that has been encrypted with the servers public key. After this, every message can be transferred between client and server in an encrypted way.

2.3. Authentication via Password

To let a user log in on the system, the user's authenticity has to be proven. To achieve that, [Linux](#) provides [PAM](#). It provides a clean separation between a program and the sensitive part of the authentication. How the server solves this situation is up to the implementation. It is also possible to rely on the [login\(1\)](#) ([2012](#)) command, which also uses [PAM](#) in the background.

2.4. User Data Querying

To spawn the default [shell](#) of a user, the path to said [shell](#) is required. This and other information can be found inside the `/etc/passwd` file. The file holds information about a user's [login-shell](#), password hash, groups, user information, [User ID \(UID\)](#), [Group ID \(GID\)](#) and home directory. To spawn a [shell](#), the first and the last of these pieces of information are substantial to successfully log in a user. To query such data, [Unix](#) offers [getpwnam\(3\)/getpwuid\(3\)](#) ([2019](#)), which does not just read the `passwd` file, but also draws information about users from other sources like [Network Information Service \(NIS\)](#) or [Lightweight Directory Access Protocol \(LDAP\)](#).

2.5. Starting the Shell

A [shell](#) requires a multitude of information before it can function normally:

- The name of the user and the host.
- The `TERM` variable to allow for the use of [ncurses\(3X\)](#) ([2019](#)) dependent applications.
- The size of the [teletype \(tty\)](#) window (including updates to the window's size).
- It has to be the session leader.

The user information depends on the user that starts the shell or more precisely: The [UID](#) and [GID](#), with which it has been started. Note that this is not the same as the display name of the user in the `USER` variable. The host name is the same as the `HOSTNAME` environment variable of the parent process.

There is also to note that for the usage of the correct terminal sequences, the `TERM` environment variable has to be known to the [shell](#). This is needed to use the correct escape sequences for applications that use [ncurses\(3X\)](#) ([2019](#)).

A [tty](#) or terminal emulator sends a [window change signal \(SIGWINCH\)](#) to its child process to notify a change in its window size. If this is not forwarded, then the width and size the shell assumes, might collide with the actual values and lead to overflowing lines and unused space.

To take full control of the terminal, the [shell](#) has to be the session leader. This can be achieved by letting it set the [Session ID \(SID\)](#). If it is not set, the [shell](#) will print some errors regarding the `ioctl` device.

2.6. Login Accounting

When a user logs into a [Unix](#) machine, a session and a time stamp get created.

For this, the two files `/var/run/utmp` and `/var/log/wtmp` provide the appropriate storage. The `utmpx` [API](#) offers appropriate functions. The `utmpx` struct represents a single entry in those files:

```

1  #define _GNU_SOURCE
2  /* Without _GNU_SOURCE the two field names below are prepended by "__" */
3  struct exit_status {
4      short e_termination; /* Process termination status (signal) */
5      short e_exit; /* Process exit status */
6  };
7  #define __UT_LINESIZE 32
8  #define __UT_NAMESIZE 32
9  #define __UT_HOSTSIZE 256
10 struct utmpx {
11     short ut_type; /* Type of record */
12     pid_t ut_pid; /* PID of login process */
13     char ut_line[__UT_LINESIZE]; /* Terminal device name */
14     char ut_id[4]; /* Suffix from terminal name, or ID field from
15                     inittab(5) */
16     char ut_user[__UT_NAMESIZE]; /* Username */
17     char ut_host[__UT_HOSTSIZE]; /* Hostname for remote login, or kernel
18                                     version for run-level messages */
19     struct exit_status ut_exit; /* Exit status of process marked as
20                                     DEAD_PROCESS (not filled in by init(8) on Linux) */
21     long ut_session; /* Session ID */
22     struct timeval ut_tv; /* Time when entry was made */
23     int32_t ut_addr_v6[4]; /* IP address of remote host (IPv4 address uses
24                             just ut_addr_v6[0], with other elements set to 0) */
25     char __unused[20]; /* Reserved for future use */
26 };

```

Listing 2.1: Definition of the `utmpx` structure ([Kerrisk 2010](#), p.819)

On [login](#), a record has to be written to the `utmp` file to indicate that the user logged in. If there is already a record for the active [terminal](#), then the entry has to be updated, otherwise a new entry has to be appended. A call to `pututxline(3)` ([2017](#)) should suffice in performing these steps properly. The application has to set the `ut_type` field of the `utmpx` struct to `USER_PROCESS` to mark a user [login](#).

Similarly to the `utmp` update after [login](#), the application has to report to the `utmpx` [API](#) that the session ended. This procedure consists of almost the same actions as the one after logging in, with exception of `ut_user` being zeroed out and `ut_type` being set to `DEAD_PROCESS` ([Kerrisk 2010](#), p.828).

A program can also query the `utmp` file with the according `get*` methods:

```

1  #include <utmp.h>
2
3  struct utmp *getutent(void);
4  struct utmp *getutid(const struct utmp *ut);
5  struct utmp *getutline(const struct utmp *ut);
6
7  struct utmp *pututline(const struct utmp *ut);

```

Listing 2.2: `utmpx` [API](#) functions

It is worth mentioning that the [Linux](#) command `login(1)` ([2012](#)) also uses `PAM` to authenticate a user.

2.6.1. Login Accounting with PAM

PAM supports binding `login` accounting to its own session functions `pam_open_session(3)` (2016) and `pam_close_session(3)` (2016):

```

1 #include <security/pam_appl.h>
2
3 int pam_open_session(pam_handle_t *pamh, int flags);
4 int pam_close_session(pam_handle_t *pamh, int flags);
  
```

Listing 2.3: PAM session management

2.7. Privilege Separation

Allowing a user to log in as any user requires the authenticating process to have `root` rights. However: After performing the authentication procedure and after successfully finishing the latter, a `shell` will be spawned for the user to interact with. This implies that without dropping the privilege of the process down to the appropriate privileges of the logged in user, the `shell` would run with `root` rights, which is a privilege escalation and not permissible.

Linux provides `setuid(2)` (2019) and `setgid(2)` (2019) to set the owning user/group of a process and therefore also changing it's permissions. Unless the user is `root`, the process' owner cannot be changed.

In comparison: SSH chains the authentication of a user between an unprivileged child process that forwards the information between the two processes (see figure 2.2) - thereby putting a privilege separation in place.

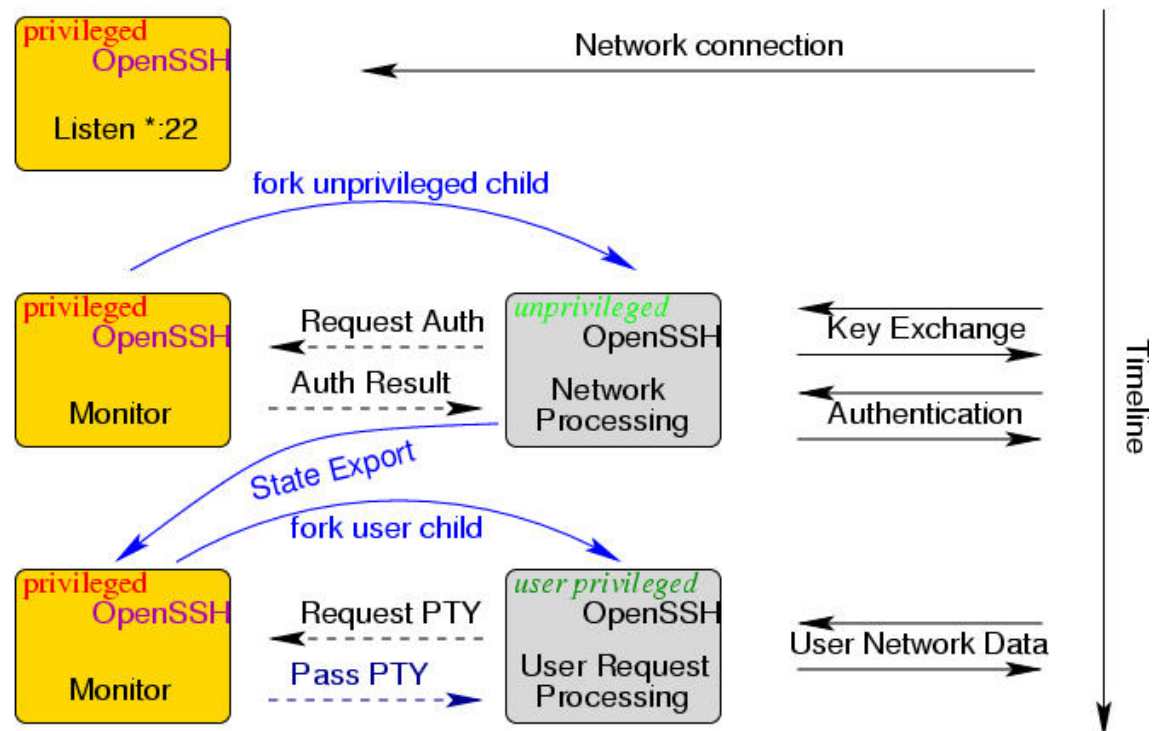


Figure 2.2.: Privilege separation in SSH (Provos 2003)

3. Implementation

The implementation of the project differs from the envisioned design. This new version of the server has three executables:

- The first is the client-side application, that is used to connect to a server.
- One is `goshd`, a `daemon` (also refer to 3.10), which only has one function: To listen for incoming connections and then to start a child process of the second executable.
- This last executable is the host (`goshh`) and it is responsible for handling the new connection.

Now, the work flow of the solution looks as follows (note that unchanged entries are gray):

1. The server listens for incoming connections.
2. The client dials server.
3. The server executes `goshh` as a child process (referred to as “host” in subsequent steps).
4. The host handles the established connection.
5. The host sets up a secure connection (see 3.2).
6. The host requests necessary environment variables from the client.
7. The client responds with the values of said variables.
8. The host notifies the client to prepare itself for the session.
9. the client sets the appropriate `tty` mode (see 3.8).
10. The host initiates a `login` procedure
 - The host finds a public key that belongs to the client and is authorized for `login` on the specific user account. The host now initiates authentication via keys (see 3.4):
 - a) The host encrypts a random secret with it and then sends that to the client.
 - b) The client decrypts the secret and sends the answer back.
 - c) The host checks whether the received answer matches the original secret:
 - If it matches, the authentication has succeeded.
 - i. The host gathers essential data about the logged in user (see 3.5).
 - ii. The host spawns the users `login shell` (see 2.5 and 3.6) with the credentials of the logged in user and forwards all traffic between `shell` and client.
 - If it fails, the authentication via keys has failed and instead, the authentication via password gets initiated.
 - The host does not find a public key that matches the client or the authentication with keys did not succeed.
 - a) The host starts an instance of `login(1)` (2012), which asks the client to authenticate itself (see 3.3).
 - b) After successful `login`, `login(1)` (2012) drops privileges, takes post-`login` actions and starts the user’s `login shell`.
11. The client uses the `shell` on remote machine.

12. The session ends in particular ways:

- The client terminates session.
 - a) The child process of the host terminates.
 - b) The host terminates.
- The client dies.
 - a) The connection dies.
 - b) The **shell** receives EOF and terminates.
 - c) The host terminates.
- The host dies or receives a **interrupt signal (SIGINT)**.
 - a) The connection dies.
 - b) The client receives EOF and terminates.
- The server dies or receives a **SIGINT**.
 - a) The server sends a **SIGINT** to all the active hosts.
 - b) The host terminates.
 - c) The connection dies.
 - d) The client receives EOF and terminates.

A new sequence diagram was created after the rough finishing of the project to display it's new work flow (see figure 3.1).

In the process of implementing this project, several problems arose that had to be addressed. The reason being that some envisioned features or mechanisms could not be implemented as originally thought. This is also the reason why the work flow described in the beginning of this chapter differs from that of the design in chapter 2.

3.1. Problems

3.1.1. Forking

To handle new established connections, it was deemed important to **fork(2)** (2017) the process, as this duplicated the current process' memory and returns the **Process ID (PID)**: 0 for the child process and a number greater than 0 for the parent to have the **PID** of the child.

In theory, this should have enabled the program to use **Go**'s standard library capabilities to handle connections. However, there were several problems with this approach, which break the privilege separation design discussed in section 2.7:

Forking not supported

The **Go** standard library does not support the classical **C**-like forking. According to Google, **Go** does not have such mechanics, as **Go** was designed with **Go**-routines in mind instead. **Go**-routines however do not support privilege dropping. Instead, it only has a `syscall.ForkExec` method, which is documented as:

Combination of **fork(2)** (2017) and **exec(3)** (2019), careful to be thread safe.

But since it uses **exec(3)** (2019) as well, it is the same as calling arbitrary binaries/scripts with the `exec.Cmd` function.

However: **Go** has a feature called **CGo**, which allows programs to call and interact with native **C**-routines. This opens up the possibility of using **fork(2)** (2017).

Forking breaks Go objects

Forking with the functionality of `CGo` does not solve the problem either. The reason is that after forking there are two programs with a `net.Conn` object. This led to both connection objects being corrupted and turning unusable. Therefore, it was necessary to abandon the clean solution of forking and instead creating a new executable that can handle new connections by its own.

Sharing Data with Child

The question then was: How can a process instantiate a child process and hand over all resources to it necessary for handling the new connection?

Since they are two separate processes now, they do not share any memory anymore. Hence the parent has to give the child the information about the connection via arguments. The most direct way to deal with this is to use `file descriptors (fds)`, which can be passed as integer arguments to the child.

Go Connection Cannot Be Transferred

Getting a `net.Conn` interface from a `fd` is supported in `Go` via:

```
1 fd := uintptr(0) // Dummy fd for context
2 conn, err := net.FileConn(os.NewFile(fd, "conn"))
3 if err != nil {
4     panic(err.String())
5 }
```

Listing 3.1: Getting a `net.Conn` interface from a `fd`

Getting the `fd` from a connection is also possible:

```
1 file, err := conn.(*net.TCPConn).File()
2 if err != nil {
3     panic(err.String())
4 }
5 fd := file.Fd()
```

Listing 3.2: Getting the `fd` from a `net.Conn` object

However: Creating a connection with the high-level `API` of `Go` and handing over the `fd` to the child to derive a `net.Conn` object from, fails.

This had some implications for the project: The listener on the server could not be created with high-level functions like:

```
1 ln, err := net.Listen("tcp", ":8080")
2 if err != nil {
3     // handle error
4 }
5 for {
6     conn, err := ln.Accept()
7     if err != nil {
8         // handle error
9     }
10    go handleConnection(conn)
11 }
```

Listing 3.3: `Go`'s high level `API` for listener

Instead the project had to rely on the low-level `socket`. The `x-package` `sys/unix` provides the necessary wrapper functions, which can be used instead. The obvious drawback being having to rely on an `x-package` which is subject to change and not being able to use the higher-level methods which **are** part of the standard library.

3.1.2. Privilege Dropping

To have a sensible privilege separation, the host process should drop the `root` privileges to the privileges of the logged in user. But after performing the `login` procedure, setting the `UID` or `GID` both resulted in an error about insufficient permissions. This problem could not be solved in the course of this project. Instead, it was relied upon spawning the `shell` already with set `UID` and `GID` values to ensure appropriate permissions. This worked out well, as looking at the process in a process monitor like `htop` showed that the spawned `shell` does have the right `UID` and `GID`.

3.1.3. Terminal Process Group

A `login shell` such as `bash` expects a properly set up `fd` with `ioctl(2)` (2017). This can be done by allowing the process to set `c tty`. However: Setting the corresponding flag when starting the process via `Go` fails with an error. The solution does not contain a solution for this problem. This shows when starting a `login shell` like `bash`, which prints the following:

```
1 bash: cannot set terminal process group (-1): Inappropriate ioctl for device
2 bash: no job control for this shell
```

Listing 3.4: `bash` error when starting `shell` without the `setc tty` flag

3.2. Secure Connection

As described in 2.2, `TLS` was used to secure the connection between the client and the server. For this, the `Go`-package `crypto/tls` has been used. This package requires both client and server to extend the `GODEBUG` variable with a flag to activate support for `TLS 1.3`:

```
1 func init() {
2     os.Setenv("GODEBUG", os.Getenv("GODEBUG")+",tls13=1")
3 }
```

Listing 3.5: Activating `TLS 1.3` in `Go`

The server-side of the solution has to have a certificate, which has to be generated newly when actually installing the solution on a machine. For testing purposes, a self-signed certificate was created using `openssl(1)` (n.d.):

```
1 openssl req -newkey rsa:2048 -nodes -keyout key.pem -x509 -days 365 -out
   certificate.pem
```

Listing 3.6: Generating a self-signed certificate and private key

The files `key.pem` and `certificate.pem` can be found inside the `test` folder in the project's root folder.

The current implementation is set to skip verification of insecure certificates. Handling such a case has been decided to be handled in future improvements. The security aspects of this implementation are discussed in subsection 5.3.1.

3.3. Authentication via Password

There are no official packages in `Go` that provide a wrapper to `PAM` (refer to 2.3). In earlier stages of the project, a wrapper from Steinert (2015) has been used. However, `login` using `PAM` failed on one of the test environments (namely Arch Linux). After some trial and error, switching to `login(1)` (2012) solved said issue, even though this command also uses `PAM`. This came with additional desired features like the `login-accounting` and other post-`login` actions (see 2.6 and 2.6.1), which the command already covers. After roughly a month of switching to `login(1)` (2012), the `login` procedure failed again on said environment. This occurred in a later stage of the project and it was deemed too time consuming, returning to the old implementation, which still lacked `login`

accounting. Fortunately this only affected one of the environments: The other environments (i.e. [Windows Subsystem for Linux \(WSL\)](#)) were still functioning.

3.4. Authentication via Keys

A user can also authenticate itself without a password but instead using public key cryptography. For this, a user has to create a key pair consisting of a private and a public key. Before the actual authentication, a hashed version of the public key has to be stored on the server.

When starting an authentication, the user sends the public key to the server, which compares its hash with the stored keys that are deemed permissible for authentication. If it matches, the server encrypts a random secret (with high entropy) with the public key to the user. The user decrypts the message with the private key and sends it back to the server. If the returned secret matches the original secret, the user proved that it is the legitimate owner of the public key.

This authentication is sufficiently secure from third parties which do not have access to the private key of the user, as it can prove the authenticity of a user.

This mechanism has been implemented but with some alterations.

The first change was to store the public key in a plain text format, which was done out of convenience and can be changed in the future. Not storing the hashed public key does not pose any security threat, so fixing this deviation was deemed of low priority.

The second and last alteration was to have the structure of authorized keys on the server be stored exclusively in the `root` user's home directory. This is a sub-optimal approach, as this requires users to store their public keys in the super user's directory, which can lead to mistakes. This would be fixed in future works.

To test the application, a test user was created and given a key pair which was created as follows using `openssl(1)` (n.d.):

```
1 openssl genpkey -out client.pem -algorithm rsa -pkeyopt rsa_keygen_bits:2048
2 openssl rsa -in client.pem -out client.pub -pubout
```

Listing 3.7: Generating a key pair for the client

3.5. User Data Querying

Querying user data (see 2.4) is *partially* supported by the [Go](#) standard library. It can give all the information about a user on the system but its [login shell](#). Therefore it was necessary to obtain said information in a different manner. One way to get the [login shell](#) of a user is to read in the `passwd` file of the system. This originally has been implemented by using a private [Go](#) package that parsed the `passwd` file. This was deemed incomplete, as a system can acknowledge users that are not listed in said file. An example of this is [NIS](#) or [LDAP](#), which would not be covered by solely relying on the `passwd` file.

In later stages, this has been corrected by switching to using [CGo](#) and calling the [Unix API](#) routines `getpwnam(3)/getpwuid(3)` (2019) as described in section 2.4. The reason for using [CGo](#) is that the `sys/unix` package does not have a wrapper function for `getpwnam(3)/getpwuid(3)` (2019).

```
1 #include <sys/types.h>
2 #include <pwd.h>
3
4 struct passwd {
5     char *pw_name; /* username */
6     char *pw_passwd; /* user password */
7     uid_t pw_uid; /* user ID */
8     gid_t pw_gid; /* group ID */
9     char *pw_gecos; /* user information */
10    char *pw_dir; /* home directory */
11    char *pw_shell; /* shell program */
12 }
```

```
12 };
13
14 struct passwd *getpwnam(const char *name);
15 struct passwd *getpwuid(uid_t uid);
```

Listing 3.8: Definition of passwd and [getpwnam\(3\)/getpwuid\(3\)](#) (2019)

3.6. Starting the Shell

In [Go](#), a process can be started with special settings. These allow for setting the [UID](#), [GID](#) and even letting the process set itself as the session leader.

```
1 cmd := exec.Command("/bin/bash", "--login")
2 cmd.SysProcAttr = &syscall.SysProcAttr{
3     Setsid: true,
4     Credential: &syscall.Credential{
5         Uid: pwd.Uid,
6         Gid: pwd.Gid,
7     },
8 }
9 cmd.Env = []string{"TERM=xterm-256color"}
```

Listing 3.9: Starting a process in [Go](#)

The transfer of the [SIGWINCH](#) has not been realized in this project because of the lack of time in the later stages of the project.

3.7. Login Accounting

The current implementation either outsources the post-login actions described in [2.6](#) to [login\(1\)](#) (2012)(see [3.3](#)) or omits them completely, as is the case in [3.4](#). The implementation of this feature is not part of the current project.

3.8. Terminal Mode

The client expects to send all its input directly to the [shell](#) without prior interpretation. To achieve this, the client-side [terminal](#) has to be set from the default cooked- into the raw-mode([Kerrisk 2010](#), p.1309). This sends each key stroke to the server-side [shell](#) as is. After the termination of the session, the original cooked-mode has to be restored to ensure operation as per usual.

Setting the [terminal](#) mode as described above in [Go](#) can be achieved with the functionality of the [x-package](#) “[golang.org/x/crypto/ssh/terminal](#)”:

```
1 import "golang.org/x/crypto/ssh/terminal"
2 //...
3 oldState, err := terminal.MakeRaw(0)
4 if err != nil {
5     panic(err)
6 }
7 defer terminal.Restore(0, oldState)
```

Listing 3.10: Setting the [terminal](#) mode in [Go](#)

3.9. Pseudoterminal

ptys are an **Inter-Process-Communication (IPC)** mechanism that help solving the problem of how two remote programs can communicate as if they were directly connected via a **terminal** (see figure 3.2).

A **shell** expects to be connected to a full fledged **tty**(or a **tty** emulator). To check whether it is inside such an environment, it uses **isatty(3)** (2019) on the **fd** it is connected to. Therefore the **fd** it is connected to should behave like a real **terminal**. This is where **ptys** come into play.

To do this, two files are created: The **pseudoterminal master (ptm)** and the corresponding **pseudoterminal slave (pts)**(see figure 3.3). **Linux** provides a **pty** generator at **/dev/ptmx**, which creates a **ptm** and a **pts**. This works by simply opening the **/dev/ptmx** file using **posix_openpt(3)** (2017). After this, a program has to grant the **pts** file ownership and permissions with **grantpt(3)** (2017), unlock it with **unlockpt(3)** (2017) and retrieve its file name with **ptsname(3)** (2017):

```
1 #define _XOPEN_SOURCE 500
2 #include <stdlib.h>
3
4 int posix_openpt(int flags);
5 int grantpt(int fd);
6 int unlockpt(int fd);
7 char *ptsname(int fd);
```

Listing 3.11: **pty** related **Linux API** functions

With the **pty** properly set up, a child bound to the **pts** side will assume it is connected to a **terminal** and also behave as such. This mechanism is key to running a **shell** and forwarding all traffic between a remote client and the **shell**. In case of **SSH** and the objective of this thesis, the layout looks like in figure 3.4.

Go does have a wrapper for **ptys** in an official package, but the code is inside an internal package, called **os/signal/internal/pty**, which uses **CGo** to call the required **API** routines described in 3.9:

```
1 // Open returns a master pty and the name of the linked slave tty.
2 func Open() (master *os.File, slave string, err error) {
3     m, err := C.posix_openpt(C.O_RDWR)
4     if err != nil {
5         return nil, "", ptyError("posix_openpt", err)
6     }
7     if _, err := C.grantpt(m); err != nil {
8         C.close(m)
9         return nil, "", ptyError("grantpt", err)
10    }
11    if _, err := C.unlockpt(m); err != nil {
12        C.close(m)
13        return nil, "", ptyError("unlockpt", err)
14    }
15    slave = C.GoString(C.ptsname(m))
16    return os.NewFile(uintptr(m), "pty-master"), slave, nil
17 }
```

Listing 3.12: **Go's pty** wrapper

For the project, this code has been altered to better fit the flow of the application.

3.10. Service Hosting

The **goshd** program has to run on a server to accept incoming requests. For ease of use, the program should be able to be run as a **daemon**. On **Linux**, this can be achieved using the service manager **systemd(1)** (n.d.), for which a unit file has been created, which is located inside the

init directory of the project. After deployment of the software, the service can be controlled using *systemd(1)* (n.d.) commands:

```
1 # Setup of goshd
2 sudo systemctl enable goshd.service
3 sudo systemctl start goshd.service
4 sudo systemctl info goshd.service
5
6 # Breakdown of goshd
7 sudo systemctl stop goshd.service
8 sudo systemctl disable goshd.service
```

Listing 3.13: goshd service control

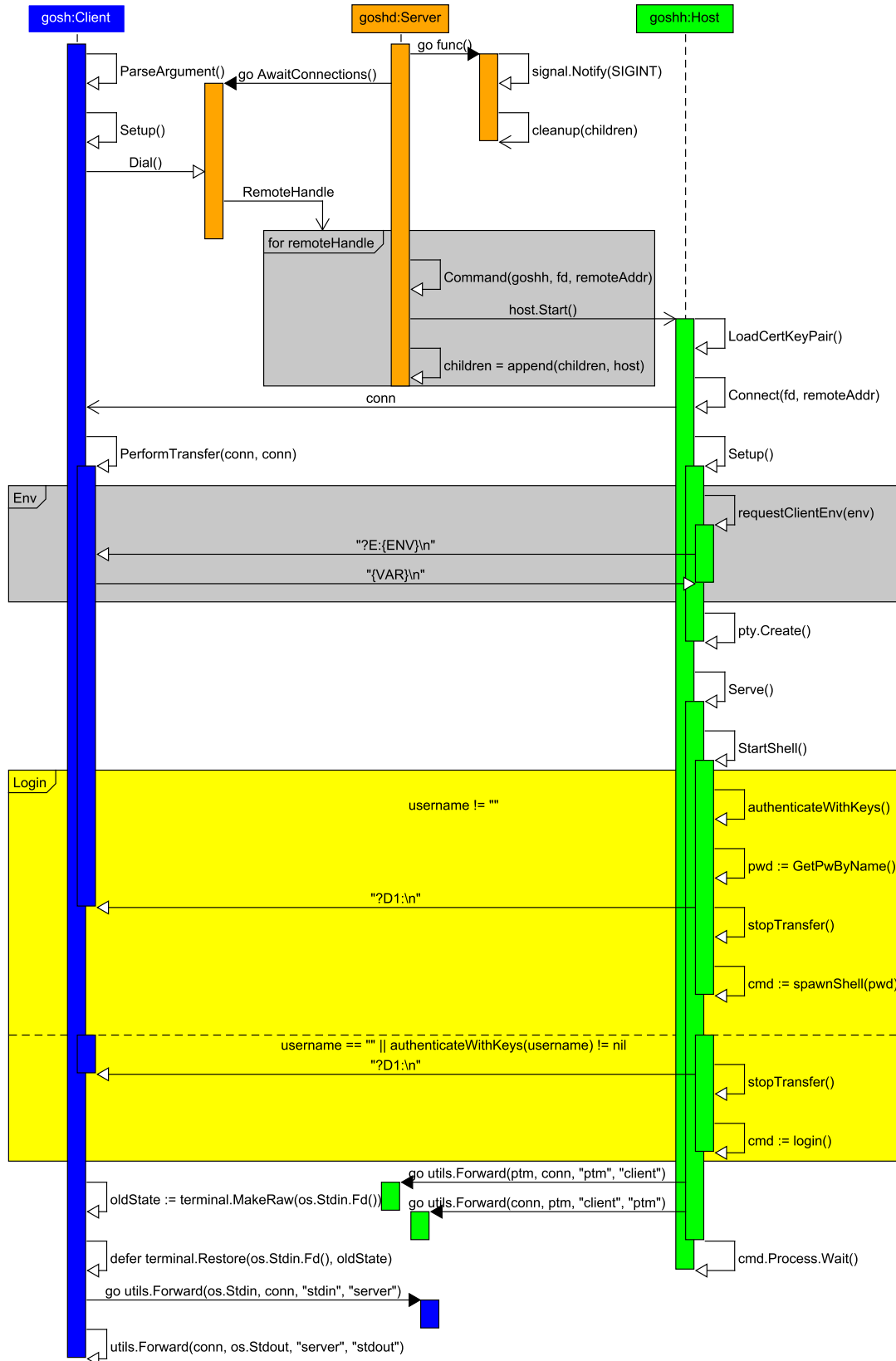


Figure 3.1.: Sequence diagram of current implementation.

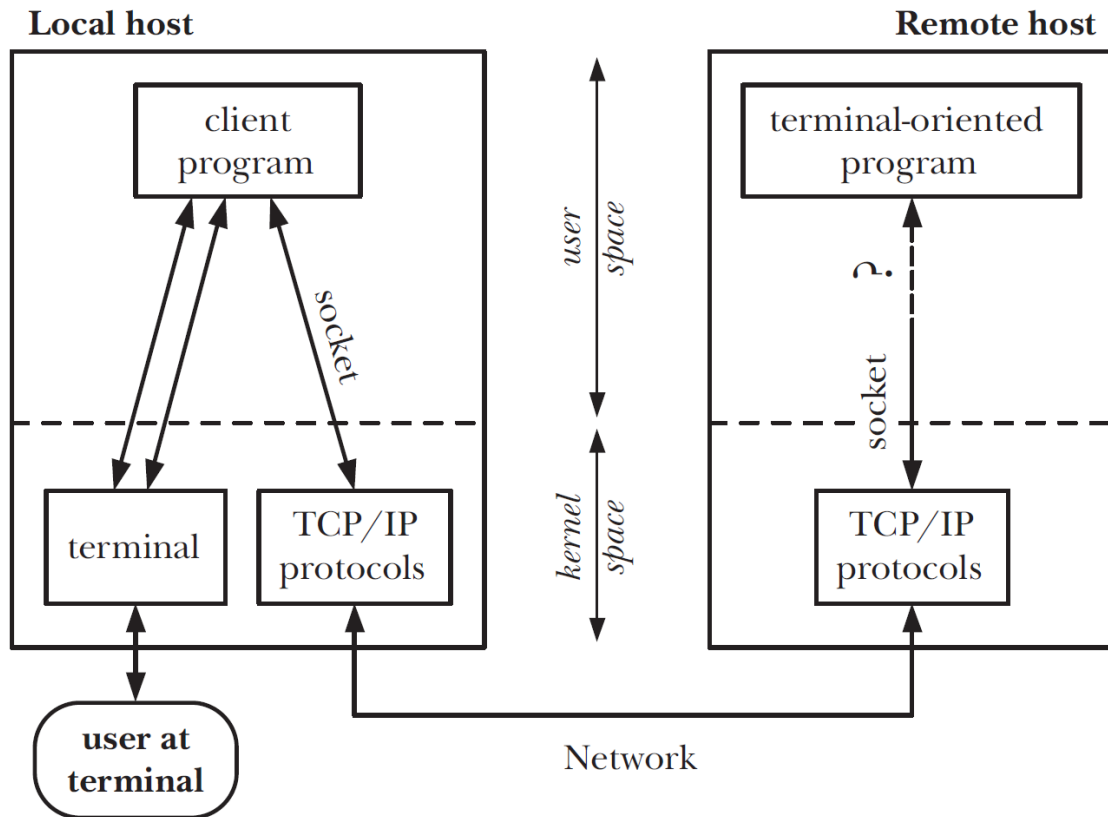


Figure 3.2.: How to operate a **ty**-oriented program over a network? (Kerrisk 2010, p.1376)

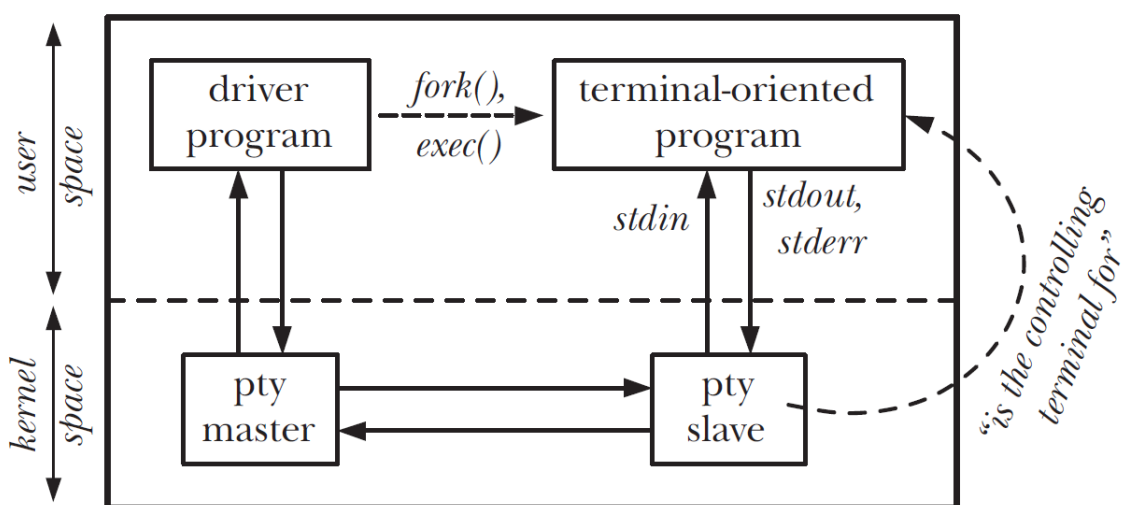


Figure 3.3.: Two programs communicating via a **pty** (Kerrisk 2010, p.1377)

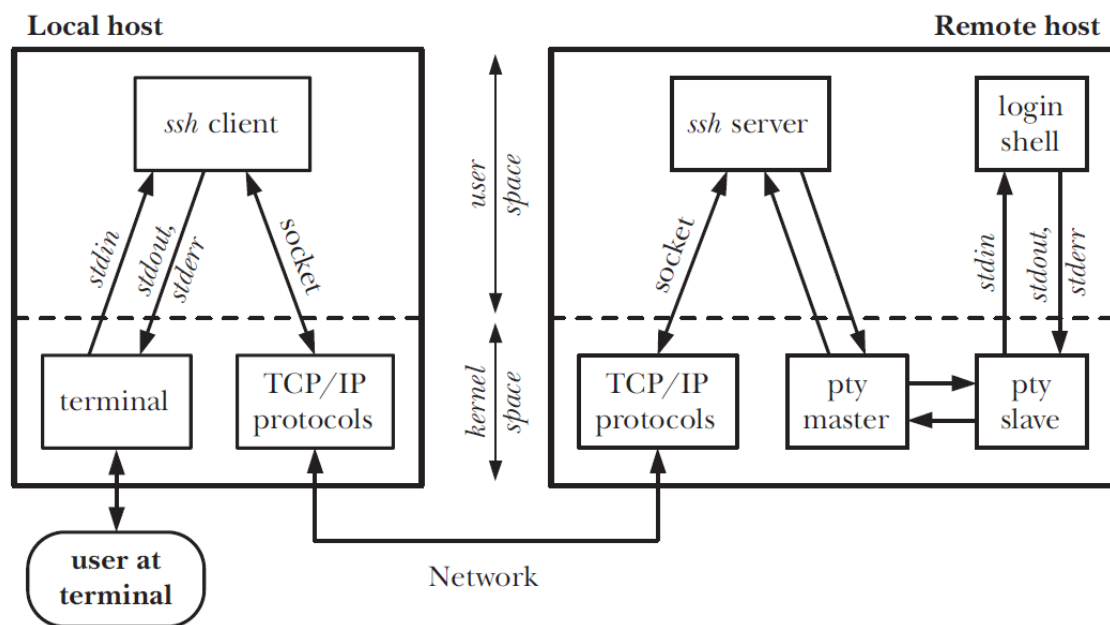


Figure 3.4.: How ssh uses a pty (Kerrisk 2010, p.1378)

- `--cert`: Sets the path of the server certificate (defaults to “`/etc/gosh/certificate.pem`”).
- `--key`: Sets the path of the server key file (defaults to “`/etc/gosh/key.pem`”).
- `goshh` is a host that handles a new connection. This binary is only executed by the server and takes the following arguments:
 - `--help`: Displays the help text.
 - `--conf`: Sets the path where the configuration file is stored (defaults to “`/etc/gosh`”).
 - `--auth`: Sets the path where the key pair is stored (defaults to “`~/.gosh`”).
 - `--cert`: Sets the path of the server certificate (defaults to “`/etc/gosh/certificate.pem`”).
 - `--key`: Sets the path of the server key file (defaults to “`/etc/gosh/key.pem`”).
 - `--remote`: Sets the address of the peer (defaults to “`localhost:2222`”).
 - `--fd`: Provides the `fd` of the connection the host has to handle.

To configure the applications, the client side has a configuration file called `gosh_config.toml`. The server's configuration file is called `goshd_config.toml` respectively, where both `goshd` and `goshh` refer to the same file. In it, the user can specify the following parameters and a few more:

- The default port to communicate over (2222).
- The default log level (info).
- The location of the authorized keys (`~/.gosh`).
- The maximum amount of allowed sessions (server-side only) (10).

The project allows a user to connect to a server and enter an interactive session with a remote user's `login shell` (see 3.3, 3.4, 3.7, 3.8 and 3.9). On the server side an appropriate privilege separation is performed upon user `login` (see 2.7 and 3.1). The communication is secured in the beginning of the transaction using `TLS` (see 3.2).

To test the applications, the user can execute the binaries from within the root folder of the project on both the client and the server (they can be the same machine, if so desired). To do this, the user can make use of the `CLI` flags described above to redirect the dependencies to the test files in the project:

```
1 # Setup environment: Create test user, its home directory and set the login
  shell.
2 ./scripts/setup.sh
3
4 # Start the server
5 sudo goshd --conf configs --auth test --cert test/certificate.pem --key
  test/key.pem
6
7 # Start the client
8 gosh --conf configs --auth test
9 # Or
10 gosh --conf configs --auth test 192.168.0.100
11 # Or
12 gosh --conf configs --auth test test@localhost
```

Listing 4.1: Running the applications for test purposes

To change the log level, giving the application the environment variable `LOG_LEVEL` set to the desired log level (i.e. “trace”, “debug”, etc.) should suffice.

5. Discussion And Prospects

As described in chapter 4, the developed solution is capable of providing a TLS secured channel to an interactive session of a remote user's shell, of which the privilege has been dropped to the appropriate level for said user. It can be said that with this, the main goal of this thesis could be achieved. In the following list, an overview of the solutions to the official tasks are given and explained:

- **Design and implement a client-server protocol that can manage interactive sessions**

This has been completed in chapters 2 and 3 respectively. Both chapters give an overview and then list certain points of interest to be highlighted.

- **Design and implement a privilege-separation architecture on the server side that allows safe dropping of privileges once a client establishes a connection**

The design has been discussed in section 2.7 and the implementation and problems encountered displayed in subsection 3.1.2 respectively. Despite the encountered problems, a privilege separation was still realized.

As described in section 1.4, the following tasks are required for a passing grade:

- **An introduction to the problem and why the envisaged solution will solve it**

This has been discussed in chapter 1.

- **A survey of related work in the area**

Related works are listed and discussed in chapter 1 as well (specifically in the sections 1.1, 1.2 and 1.3).

- **A detailed design of the solution**

The entire design can be found in chapter 2. The implementation differs from the original design however, which is also discussed in chapter 3.

- **An evaluation of the performance of the implemented solution**

An elaboration to this task can be found at subsection 5.1.

- **A privilege-separation architecture**

This is the same as the first task and can be considered solved as of section 2.7 with remark to subsection 3.1.2.

In the following lists, there are additional tasks described in section 1.4:

- **A comparison of all the related work with the envisaged solution, outlining why the envisaged solution is better**

A comparison to the surveyed solutions can be found in section 5.2.

- **A detailed analysis of the security of the solution, including possible attacks and defenses**

This task has been cleared in section 5.3.

- **Use of TLS as the transport layer**

This topic can also be considered solved as of sections 2.2 and 3.2.

- **A proof-of-concept client that can handle interactive sessions**

A usable client has been developed and can be used after compilation of the source code. An overview of the developed applications can be found in chapter 4.

- **A proof-of-concept client that works as a transport for rsync**

This is a task that has not been solved in this thesis but its addition is discussed in subsection 5.4.2.

5.1. Performance

The `pty` is an interface that needs data to be forwarded to and received from. This holds true for both sides (`ptm` and the `pts`) and therefore needs a mechanism to transfer said data. On the slave side, the `shell` itself already seamlessly communicates with the device via its standard pipes. A similar concept is applied for the client and its connection to the server, where the application communicates via the connection with the server. There, the client uses a `Go`-routine to asynchronously forward the standard input of the program to the connection.

The master side has to transfer the input from the remote client to the `ptm` and the output from the latter to the former as well. For this, the server uses two `Go`-routines that asynchronously forward the data from each stream to the other and waits until the shell process terminates.

Both these transfers are performed using `Go`'s `WriteTo` method:

```
1 // WriteTo implements io.WriterTo.
2 // This may make multiple calls to the Read method of the underlying Reader.
3 // If the underlying reader supports the WriteTo method,
4 // this calls the underlying WriteTo without buffering.
5 func (b *Reader) WriteTo(w io.Writer) (n int64, err error)
```

Listing 5.1: `WriteTo` method of `Go`

This mechanism helps keeping the memory usage low.

As for speed, the application has been tested on speed of the connection. For this, the client was changed to sending a stream of zeroes from `/dev/zero` to the connection once the latter has been established. The server-side on the other hand still spawns a `host`, but that forwards the entire data received from the connection to its `stdout`. To measure the throughput, the server's output gets piped into `pipeviewer` (`pv`), with the following arguments:

- `-t`: show elapsed time
- `-r`: show data transfer rate counter
- `-a`: show data transfer average rate counter
- `-b`: show number of bytes transferred
- `-W`: display nothing until first byte transferred

The data stream gets then redirected to `/dev/null`.

```
1 # Set the log level to error to prevent sending bytes to pv before the client
   # sends the stream
2 sudo LOG_LEVEL=error goshd | pv -rabtW > /dev/null
```

Listing 5.2: Throughput measurement

Two variations of the test were performed: One using the blank `Go` connection `net.Conn` and the other one having a `TLS` layer in between. The equipment for this test consisted of a native Arch `Linux` laptop and an Ubuntu `Linux` running inside a native Window's `WSL`.

Using this as a basic setup, several tests were performed:

- On a native Arch `Linux` laptop, the whole solution was tested over `loopback`.

- Similarly, the solution was also tested over [loopback](#) in [WSL](#).
- Connected via a switch, [WSL](#) hosted the server. The same Arch [Linux](#) laptop as in the former test connected to the server as the client.

All the runs were performed for exactly one minute. The measured speeds are as follows:

	TLS (size)	no TLS (size)
Arch Linux (loopback)	427 MiB/s (25.1 GiB)	1.15 GiB/s \simeq 1177.6 MiB/s (69.0 GiB)
WSL (loopback)	69.7 MiB/s (4.09 GiB)	116 MiB/s (6.82 GiB)
Arch Linux to WSL (ethernet)	85.1 MiB/s (4.99 GiB)	83.7 MiB/s (4.91 GiB)

Table 5.1.: Throughput measurements

The throughput in [WSL](#) never deviated by more than 8 MiB/s, whereas Arch [Linux](#) had deviations in the scale of 0.3. A reason for this was not found. The reason for the comparably low throughput over a network, however, was assumed to be the combination of the Netgear switch and a Cat 5 ethernet cable.

This demonstrates that the application is sufficiently fast to be used in most use cases. Therefore the feature of this solution to be used as a transport protocol (i.e. for [rsync\(1\)](#) (2018) - see [5.4.2](#)) is deemed viable.

5.2. Comparison

5.2.1. OpenSSH

As described in section [1.1](#), [SSH](#) uses an own protocol to secure the connection between client and server. This project differs in this point, as it uses [TLS](#) to create a secure channel (see [2.2](#) and [3.2](#)).

Both [SSH](#) and the developed solution provide means to login with a user/password combination or via keys. This feature behaves the same in both cases.

[SSH](#) has a plethora of features that are not supported by the solution, but also were not required. The reason for that is simply because the task was to create a simple alternative.

Both have privilege separation mechanisms implemented, but the developed solution was not able to drop the privilege of the host process after [login](#). This is another point where the two differ. That the solution does not drop privileges after [login](#) is also listed as a security issue in section [5.3](#).

Authentication in [SSH](#) are per default solved using [PAM](#), which can be changed in the configuration files. Although the project in earlier stages used [PAM](#) as well, it outsourced the [login](#) functionality to [login\(1\)](#) (2012), which in turn uses [PAM](#) as well.

[SSH](#) also performs [login](#) accounting in each case, whereas the solution again outsources this mechanism to [login\(1\)](#) (2012) or omits it completely (in the case of authentication via keys). This is a feature that is missing from the project (see [5.4.5](#)).

5.2.2. Telnet

In comparison to [telnet\(1\)](#) (1994), the solution has several differences, that distinguishes it from the old protocol. For one, the communication channel of [telnet\(1\)](#) (1994) is not secured. The developed solution however provides a [TLS](#) secured communication tunnel. This is the main difference between these two.

Furthermore, [telnet\(1\)](#) (1994) does not support authentication via keys, which is a feature that has been implemented in the project. Authentication via keys is also performed using [PAM](#).

Just as in the case of [SSH](#), it also does proper [login](#) accounting, which as mentioned before is missing in the solution when authenticating via keys.

5.2.3. Berkeley r-Commands

From the list of Berkeley r-commands, only *rlogin(1)* (1999) is of interest, as the other ones are used with slightly different use cases in mind. The only other one that could be compared to as well is *rsh(1)* (1999), which without a command to execute on the remote server behaves the same as *rlogin(1)* (1999).

As mentioned in section 1.2, the communication channel of *rlogin(1)* (1999) is not secured. Just as with *telnet(1)* (1994), user-password combinations are sent over a network in plain text and can be read without much difficulty. This gets prevented in the developed solution by relying on TLS.

The r-commands also do not support authentication via keys, which has been implemented in the solution. The authentication via password on the other hand is virtually the same, as both use *login(1)* (2012) to accomplish this task. This also goes for using PAM for authentication, *login* accounting and the post-*login* actions the server performs.

5.3. Security

Although the solution uses TLS as a security measure, the current state of the project still has several weaknesses. Some of them seem to be of a lower magnitude, whereas others have a higher priority to be resolved. This section lists some security issues that have been solved and others that are still unsolved.

5.3.1. Secure Connection

A secure channel is mandatory in an architecture like this. As described in sections 2.2 and 3.2, the connection has been secured using TLS (specifically TLS 1.3). This means an attacker has to deal with the security measures that TLS brings with itself.

Apart from that, the current solution only uses a server-side certificate; The client's certificate is never checked as it does not exist in the current solution. This means an attacker that knows the *login* credentials of a user on the remote system can log in in the system and deal harm to it without the server being able to distinguish the attacker machine from the expected client. Using client-side certificates could help to increase the level of difficulty of an attacker to perform this attack.

Furthermore, in the current solution, insecure server certificate errors simply get ignored. An attacker can use this to take on the role of a "man in the middle". This attack vector allows the attacker to appear to the client like it is the server the client wants to log in to and at the same time appear to the server as if it is the client trying to connect to it. This leads to two TLS connections: One from the client to the attacker and one from the attacker to the server. This means the attacker can read the entire data stream in an unencrypted way and manipulate as desired. This also leads to possible password exposure to the attacker if the client authenticates itself via password to the server. Therefore this weakness is considered a severe security issue that can easily be exploited by an attacker by simply sitting between the two connecting parties. To prevent such attacks, the certificate verification cannot be skipped.

5.3.2. Authentication

As described in sections 2.3, 3.3 and 3.4, two possible paths of authentication were realized.

Authentication via Password

Relying on *login(1)* (2012) simplified the task, as time was pressing. Using this standardized command also helps removing possible security issues in the authentication process and afterwards. However, this comes with a cost: The *login(1)* (2012) command allows a user to directly log in as root user. This cannot be prevented by the server. Ideally the server would have a default option of allowing root login set to false. There is no reason to allow any client to try to log in as root user. Especially since a logged in user can simply use su or sudo to gain root privileges. "root"

however is the only user name that can be logged in and an attacker knows exists and only needs the knowledge of one password (which *could* be a default password) to gain full system control.

Another way to prevent `root` user login despite the usage of [login\(1\)](#) (2012) could be to pre-fetch the user credentials from the client and use the `-f` flag with [login\(1\)](#) (2012) (see [5.4.3](#)).

Authentication via Keys

When authenticating a user via public key cryptography, the server looks up which users can be logged in. It does so by looking into the `root` user's home directory `/root/.gosh`. There, the directory structure defines what public key can log in as which user. In this case, Each directory has the name of a user and inside are the public keys stored in plain text with the file name of the `USER` value of the connecting client. This for one means that to allow a remote user to log in as this very user, the remote user's public key has to be named after that user's name and placed into the `root` user's home directory. Storing such files in the `root` user's home directory is a bad idea as this means that every change to this file structure needs `root` privileges. The public keys are also stored in plain text, which is not needed and can be replaced with their hashes instead as described in section [3.4](#).

5.3.3. Privilege Separation

To have a proper privilege separation (see section [2.7](#)), the host process should also drop its privileges once the [login](#) succeeded. This is desired as the forwarding of the data is still in the host's hands and as the host still runs with `root` privileges, an attacker could potentially gain `root` access if the attacker manages to gain access to that process. Such privilege escalation should be prevented but was not entirely achieved in the current solution, as described in subsection [3.1.2](#).

5.4. Prospects

Although the project created a working alternative to [SSH](#), there are several points that have to be improved, which could not be completed in the course of this thesis. These points are listed here:

5.4.1. Security Issues

Of course the primary task to solve in further works would be to solve the security issues with this project. As this project should be able to replace [SSH](#) in its core functions, the security aspects have to be comparable as well. Suggestions on how to solve the above mentioned security issues can be found in the respective subsections [5.3.1](#), [5.3.2](#) and [5.3.3](#).

5.4.2. Rsync

One of the proposed extensions was to use this project's solution as a transport layer for [rsync\(1\)](#) (2018). This could not be achieved in the course of this project, but could be done in subsequent works.

5.4.3. Prefetching of Credentials

Instead of forwarding the client's inputs to the [login](#) mechanism, the client could ask the user for user name and password beforehand and send them to the server. The server can then check for the user name to be sane (for example for preventing `root` login) and only then forward it to [login\(1\)](#) (2012). This could solve the security issue discussed in subsection [5.3.2](#).

5.4.4. Using PAM

In earlier stages of the project the authentication via password was already partially solved by using [PAM](#). For other reasons (see [3.3](#)) this was dropped in favor of [login\(1\)](#) ([2012](#)). It is suggested to return to using [PAM](#) directly and also handle [login](#) accounting (see [2.6](#)) as well.

5.4.5. Using Utmpx

Part of post-[login](#) actions are [login](#) accounting, which would be realized using [Linux](#)' [utmpx-API](#). This does happen in the case of relying on [login\(1\)](#) ([2012](#)), but not in the case of authentication via keys. The session is therefore not properly set up and should be added in future works.

5.4.6. Add Transfer of SIGWINCH

As mentioned in section [3.6](#), the [SIGWINCH](#) is not being transferred to the [shell](#). This means that the [shell](#) assumes default values for window width and height, which leads to faulty formatting. This feature should be implemented to make the interactive session more seamless.

5.4.7. Forwarding of Data Flow

As of now, the solution simply calls forwarding routines to transfer the data flow between the [shell](#) (of more specifically, the [ptm](#)) and the connection to the client. This sometimes leads to an unclear session breakdown when terminating. It is therefore desirable to implement the data forwarding with other means. Possible [Linux](#) support could be drawn from the functions [select\(2\)](#) ([2019](#)) or [poll\(2\)](#) ([2019](#)).

5.4.8. Terminal Process Group

As described in section [3.1.3](#), the [shell](#) process does not have the ability to set the terminal process group. This is a flaw in the solution that should be removed by properly allowing the process to do so.

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6.2. Glossary

Application Programming Interface

Accessible interface for developers to use external code. [1](#)

cGo [C](#) support for [Go](#). [9](#), [10](#), [12](#), [14](#)

Command Line Interface

A text based interface centered around commands to perform specific tasks. [III](#)

daemon A (possibly latent) background process on a [Unix](#) machine that handles certain events when said events occur. [8](#), [14](#), [19](#)

Diffie-Hellman key exchange A key exchange algorithm that prevents exposure to eavesdropping third parties. [5](#)

file descriptor

A file descriptor is an integer that represents the handle to a file. [10](#)

Graphical User Interface

Graphical interface for the user to visually interact with a program. [1](#)

Group ID

The unique identifier of a group represented as an integer. [5](#)

Inter-Process-Communication

Communication between processes. [14](#)

interrupt signal

A signal supported by [Unix](#) based [Operating Systems\(OSes\)](#) which signals to a process to interrupt it's work. [9](#)

Lightweight Directory Access Protocol

A protocol for querying and managing information of distributed directory information services. [5](#)

Linux A [Unix](#) based [OS](#) which uses Linus Torvalds kernel and was inspired by Minix. [II](#), [5–7](#), [11](#), [14](#), [22](#), [23](#), [26](#), [31](#), [34](#)

login The action of logging in. plural [1](#), [3–8](#), [11–13](#), [20](#), [23–26](#)

loopback A channel with only one endpoint: Sender and receiver are on the same machine. [22](#), [23](#)

Network Information Service

A service for distributing configurations like user information among a network. [5](#)

Object Oriented Programming

A programming paradigm which uses objects to model real life entities. [31](#)

Pluggable Authentication Module

Modules for user authentication. [II](#)

port A point for traffic to flow, represented by an unsigned integer of up to 2 bytes. The name was chosen as an analogy to ports for ships. [1](#), [31](#)

Process ID

The unique identifier of a process represented as an integer. [9](#)

pseudoterminal

A mechanism of [Unix](#) to allow programs to communicate as if the other was inside a [tty](#) ([pty\(7\) 2017](#)). [II](#)

pseudoterminal master

The master of a [pty](#). [14](#)

pseudoterminal slave

The slave of a [pty](#) ([pts\(4\) 2013](#)). [14](#)

Secure Shell

An client-server-application that allows remote login and interaction with a [shell](#). See [1.1](#). [II](#)

Secure Sockets Layer

Cryptographic protocol to secure the communication between two peers via symmetric cryptography. Deprecated. [5](#)

Session ID

The unique identifier of a session represented as an integer. [5](#)

shell A [CLI](#) program, that reads user input line-by-line and executes those commands. [II](#), [1–5](#), [7–9](#), [11–14](#), [20–22](#), [26](#), [31](#), [34](#)

socket A network socket is an endpoint for communication over [ports](#). [10](#)

teletype

Originally a device that could send and receive text messages. Nowadays [tty](#) refers to [terminals](#), which emulate that behaviour ([tty\(4\) 2019](#)). [5](#)

Telnet Secure

Telnet with [SSL](#) encryption. [1](#)

terminal An user interface to interact with [CLI](#) programs like [shells](#). [1](#), [6](#), [13](#), [14](#), [31](#), [34](#)

the C programming language

Low-level programming language originally invented by Dennis Ritchie. [4](#)

the C++ programming language

Descendant of [C](#) which implemented [Object Oriented Programming \(OOP\)](#). [4](#)

the Go/Golang programming language

Google's programming language. [III](#)

Transport Layer Security

Newer and recommended version of [SSL](#). [5](#)

Unix An originally free [OS](#) family called Unics from AT&T that re-imagined an older [OS](#) by the name of Multics. [II](#), [3](#), [5](#), [6](#), [12](#), [30](#), [31](#)

uptime The time a computer has been running. [2](#)

User ID

The unique identifier of a user represented as an integer. [5](#)

window change signal

A signal supported by [Unix](#) based [OSes](#) which signals to a process that the terminal window it is running in has changed size. [5](#)

Windows Subsystem for Linux

A subsystem provided by Windows that is able to run headless [Linux](#) distributions. [12](#)

x-package A [Go](#) package that is not part of the standard library and that is subject to change or even entirely disappear. plural [10](#), [13](#)

X.509 An international standard for certificates in public key infrastructures. [5](#)

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6.5. Acronym Glossary

- SIGINT** *interrupt signal* [9](#), See [interrupt signal](#)
- SIGWINCH** *window change signal* [5](#), [13](#), [26](#), See [window change signal](#)
- API** *Application Programming Interface* [1](#), [6](#), [10](#), [12](#), [14](#), [26](#), [34](#), See [Application Programming Interface](#)
- C** *the C programming language* [4](#), [9](#), [30](#), [31](#), See [the C programming language](#)
- C++** *the C++ programming language* [4](#), See [the C++ programming language](#)
- CLI** *Command Line Interface* [III](#), [20](#), [31](#), See [Command Line Interface](#)
- fd** *file descriptor* [10](#), [11](#), [14](#), [20](#), [34](#), See [file descriptor](#)
- GID** *Group ID* [5](#), [11](#), [13](#), See [Group ID](#)
- Go** *the Go/Golang programming language* [III](#), [1](#), [4](#), [9–14](#), [22](#), [30](#), [32](#), [34](#), See [the Go/Golang programming language](#)
- GUI** *Graphical User Interface* [1](#), See [Graphical User Interface](#)
- IPC** *Inter-Process-Communication* [14](#), See [Inter-Process-Communication](#)
- LDAP** *Lightweight Directory Access Protocol* [5](#), [12](#), See [Lightweight Directory Access Protocol](#)
- NIS** *Network Information Service* [5](#), [12](#), See [Network Information Service](#)
- OOP** *Object Oriented Programming* [31](#), See [Object Oriented Programming](#)
- OS** *Operating System* [30](#), [31](#)
- PAM** *Pluggable Authentication Module* [II](#), [1](#), [5](#), [7](#), [11](#), [23](#), [24](#), [26](#), [34](#), See [Pluggable Authentication Module](#)
- PID** *Process ID* [9](#), See [Process ID](#)
- ptm** *pseudoterminal master* [14](#), [22](#), [26](#), See [pseudoterminal master](#)
- pts** *pseudoterminal slave* [14](#), [22](#), See [pseudoterminal slave](#)
- pty** *pseudoterminal* [II](#), [14](#), [17](#), [18](#), [22](#), [31](#), [33](#), [34](#), See [pseudoterminal](#)
- SID** *Session ID* [5](#), See [Session ID](#)
- SSH** *Secure Shell* [II](#), [III](#), [1](#), [2](#), [7](#), [14](#), [23](#), [25](#), See [Secure Shell](#)
- SSL** *Secure Sockets Layer* [5](#), [31](#), See [Secure Sockets Layer](#)
- TELNETS** *Telnet Secure* [1](#), See [Telnet Secure](#)
- TLS** *Transport Layer Security* [5](#), [11](#), [20–24](#), [34](#), See [Transport Layer Security](#)
- tty** *teletype* [5](#), [8](#), [14](#), [17](#), [31](#), [33](#), See [teletype](#)
- UID** *User ID* [5](#), [11](#), [13](#), See [User ID](#)
- WSL** *Windows Subsystem for Linux* [12](#), [22](#), [23](#), See [Windows Subsystem for Linux](#)
- ZHAW** *Zurich University of Applied Sciences* See [Zurich University of Applied Sciences](#)

A. Appendix

A.1. Project Management

A.1.1. Official Statement of Tasks

Bachelor Thesis

Preventing Supply Chain Insecurity by Authentication on Layer 2

Stephan Neuhaus

2017-06-15

1 Introduction

The SSH protocol [RFC253, RFC6668, RFC8268, RFC8308, RFC8332] is now over twelve years old in its current form. One of the problems with SSH is its complexity, both in the initial phase when key material is exchanged, but also later, for example because the server must always decide whether to return a character that has been sent to it or not (echo).

The goal of this work is a radically simplified protocol, which in its functions is similar to SSH. (N.B. the similarity concerns the functions, not necessarily the protocol details). You develop the protocol, as well as a client and a server. You demonstrate that your software can replace SSH. For a merely passing grade, this may be done by specifying and implementing a suitable protocol. For an improved grade, you show that your replacement can handle several common use cases, among them:

- Interactive session
- Rsync with the SSH replacement as transport protocol

2 Task

To this end, this thesis will

- design and implement a client-server protocol that can manage interactive sessions
- design and implement a privilege-separation architecture on the server side that allows safe dropping of privileges once a client establishes a connection

For a passing grade (4.0), the work must contain at least the following:

- in the thesis, an introduction to the problem and why the envisaged solution will solve it;
- in the thesis, a survey of related work in the area;
- in the thesis, a detailed design of the solution;

- in the thesis, an evaluation of the performance of the implemented solution; and
- in the software, a privilege-separation architecture.

These requirements do not contain anything related to security. This is not an accident.

Incorporating the following components will improve the grade. The more components are included, the better the grade will be.

- In the related work section of the thesis, a comparison of all the related work with the envisaged solution, outlining why the envisaged solution is better;
- in the thesis, a detailed analysis of the security of the solution, including possible attacks and defenses;
- use of TLS as the transport layer;
- a proof-of-concept client that can handle interactive sessions;
- a proof-of-concept client that works as a transport for rsync;

ZHAW's School of Engineering no longer provides formal language lessons for its students as part of the curriculum. I am therefore giving notice that submitting a thesis with large amounts of orthographical or grammatical errors lead to a lower grade.

The thesis can be submitted in German or English. English is preferred, but submitting in German will not lead to a lower grade.

A.1.2. Project Plan

Revised Project Plan

March 2019

1 Introduction

The work on the final thesis is divided into several sub tasks. The individual tasks and respective time planning was defined early.

As this is a field of work, we as a team are not familiar with, we have decided to change our original plan: "Project Plan" to a revised version.

The biggest difference is that the Prototype is developed earlier but we are removing some security measures respectively moving it to an optional goal We are subdividing the following area of development:

- Technical research: The research starts with a collection of examined current solution to a secure data transference, followed by a list of pro and cons for the approaches that includes a preferred selection. Moreover, we would survey the current state development within that field.
- Conceptualising: Look for alternative solutions and compare them. Plan the development.
- Prototype: A unsecure SSH Client
- Testing: Do generalized tests, identify possible security vulnerability
- Rework of the secure shell: Modify the secure shell based on the newly discovered needs

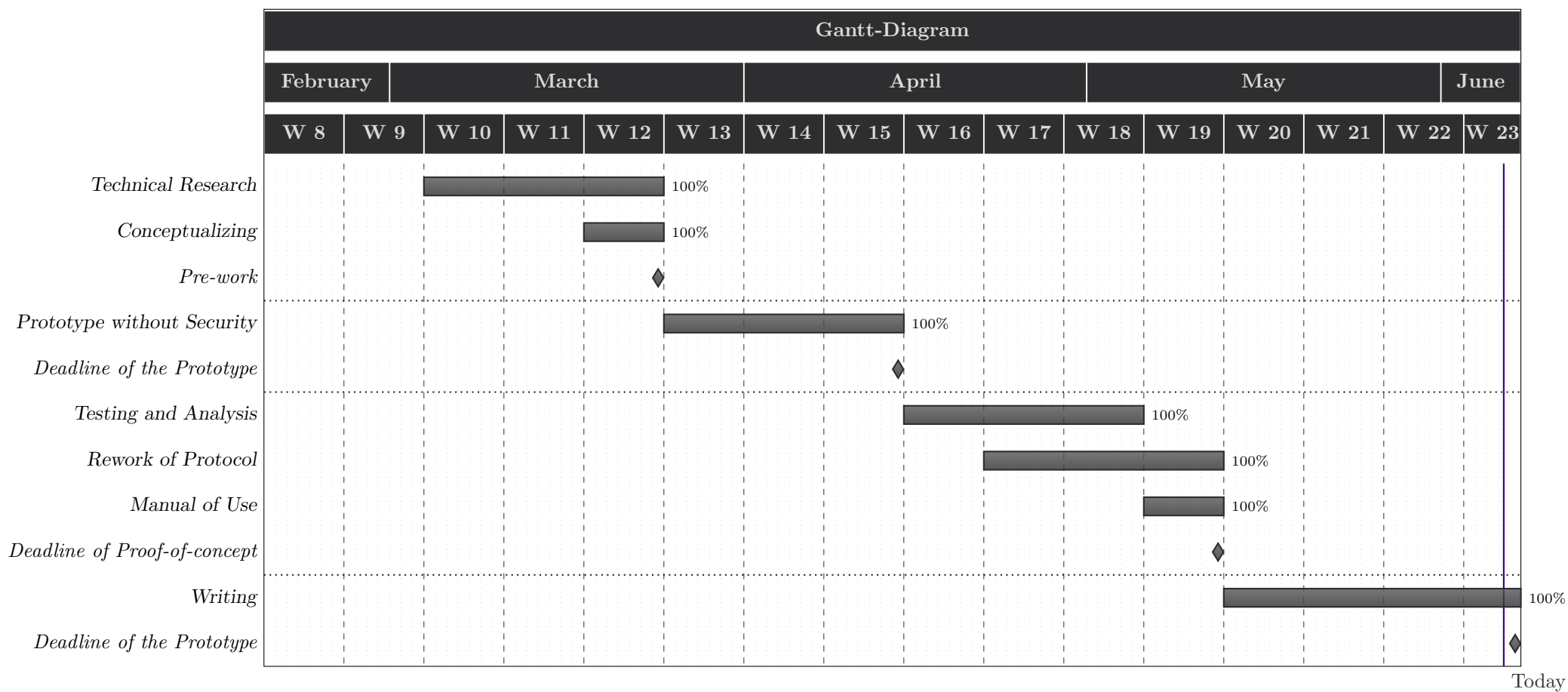
Furthermore there are specific milestone within the project process, which we would use to realign and discuss our time division.

2 Visualization

The project plan is documented in the form of a chart and is updated throughout the project. This way, deviations can be detected early and can be discussed with the supervisor and within our team.

		March				April				May				June				July
	duration	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
Technical research	3 weeks																	
Conceptualising	1 week																	
Pre-work	31.03																	
Prototype without security measures	4 weeks																	
Deadline of the Prototype	30.04																	
Testing and Analysis	3 weeks																	
Rework of current protocol	3 weeks																	
Manual of Use	1 week																	
Deadline of Proof of concept	10.05																	
Writing	2 weeks																	
Hand-in Date	28.06																	

Figure 1: Project plan



A.1.3. Meeting Minutes

The meeting minutes have a disruption in style and execution beginning from the 6th meeting. The reason for this is because in the beginning of the project, Mr Schwarz was responsible for keeping the minutes, but he opted out of the project.

Oh my Gosh - Meeting protocol

1 2nd Meeting

Participant

Bachelor thesis supervisor - Stephan Neuhausen

Bachelor student - Raphael Emberger

Bachelor student - Kevin Schwarz

Time duration of the meeting

1 hour and 15 minutes

1.1 Objectives

1. Keeping the bachelor thesis supervisor informed on the state of affairs
2. Getting an overview on how the progress relate to the scheduled progress

1.2 Summary

In this meeting the following things were achieved:

1. Decision towards ITC and UDP was achieved
2. A rough draft of a generalized process was finalized and presented to the supervisor
 - (a) Smaller misconceptions were resolved
 - (b) Fields where further research is warranted was shown
3. Reiteration of the project goal
4. Further Delimitation of the project extent.
 - (a) Smaller misconceptions were resolved

1.3 Tasks and resources

The following were left as tasks or as a research subject for the next meeting in descending priority:

1. Understanding shell forwarding
2. Researching the limitation of IOCTL - raw and device specific output
3. Pseudo terminals
4. Persistence in relation to Environment variables

1.4 Next meeting

The next meeting plan were not changed, the formerly decided weekly scheduled date still stands.

The next meeting is dated: [14 / 03 / 19] on the first floor of the Zürich location of the ZHAW within the Room 0.03.

Oh my Gosh - Meeting protocol

1 3rd Meeting

Participant

Bachelor thesis supervisor - Stephan Neuhausen

Bachelor student - Raphael Emberger

Bachelor student - Kevin Schwarz

Time duration of the meeting

1 hour

1.1 Objectives

1. Keeping the bachelor thesis supervisor informed on the state of affairs
2. Getting an overview on how the progress relate to the scheduled progress

1.2 Summary

In this meeting the following things were achieved:

1. General overview of a PTY
2. Certain foundation towards developing a non-secure secure shell were explained:
 - (a) Smaller misconceptions were resolved
 - (b) Fields where further research is warranted was shown
3. Reiteration of the project goal
4. Further Delimitation of the project extent.
 - (a) Smaller misconceptions were resolved

1.3 Tasks and resources

The following were left as tasks or as a research subject for the next meeting in descending priority:

1. Understanding shell forwarding
2. Researching the limitation of IOCTL - raw and device specific output
3. Pseudo terminals
4. Persistence in relation to Environment variables

1.4 Next meeting

The next meeting plan were not changed, the formerly decided weekly scheduled date still stands.

The next meeting is dated: [14 / 03 / 19] on the first floor of the Zürich location of the ZHAW within the Room 0.03.

Oh my Gosh - Meeting protocol

1 4th Meeting

Participant

Bachelor thesis supervisor - Stephan Neuhausen

Bachelor student - Raphael Emberger

Bachelor student - Kevin Schwarz

Time duration of the meeting

25 Minutes

1.1 Objectives

1. Keeping the bachelor thesis supervisor informed on the state of affairs
2. Gain an introduction to encryption

1.2 Summary

In this meeting the following things were achieved:

1. Introduction to Bash was given
2. Move up of the prototype deadline
3. Change of scheduled development plan:
 - (a) Decrease of the SSH scope
 - (b) Reduction of the security measures to an optional goal
 - (c)

1.3 Tasks and resources

The following were left as tasks or as a research subject for the next meeting in descending priority:

1. Researching the infrastructure of GO-order
2. Researching Bash and PTY on Windows enviroment
3. First Server - Client Demo
4. Crafting a simple process diagram for the next meeting

1.4 Next meeting

The next meeting plan were not changed, the formerly decided weekly scheduled date still stands.

The next meeting is dated: [28 / 03 / 19] on the first floor of the Zürich location of the ZHAW within the Room 0.03.

Oh my Gosh - Meeting protocol

1 5th Meeting

Participant

Bachelor thesis supervisor - Stephan Neuhausen

Bachelor student - Raphael Emberger

Bachelor student - Kevin Schwarz

Time duration of the meeting

45 Minutes

1.1 Objectives

1. Keeping the bachelor thesis supervisor informed on the state of affairs
2. Demonstrate Demo

1.2 Summary

In this meeting the following things were achieved:

1. The Prototype was tested.
2. Process Diagram was explained
3. 4 open Problems were discussed:
 - (a) PAM Struct and how they work
 - (b) Generalized Certkey location
 - (c) Use of the Prototype within Linux
 - (d) Correct pipe lining and forking

1.3 Tasks and resources

The following were left as tasks or as a research subject for the next meeting in descending priority:

1. Further testing of both Linux, Apple and Windows environment
2. Solving of Pam Struct problem
3. Further development

1.4 Next meeting

The next meeting plan were modified, the formerly decided weekly scheduled date still stands.

The next meeting is dated: [5 / 04 / 19] on the first floor of the Zürich location of the ZHAW within the Room 0.13.

Design and Implementation of an Alternative to SSH

Meeting Minutes

1 Attendees

Present: Stephan Neuhaus, Raphael Emberger

Absent: Kevin Schwarz(*illness*)

2 Initiation

The meeting took place on the *Friday, 5th of April 2019, 13:00* in *ZL0.13, Lagerstrasse 45, Zürich*. Raphael Emberger was responsible for the minutes.

3 Points of discussion

3.1 Process forking unsuccessful

Attempts on forking a sub-process were unsuccessful. The reason for this was that the standard library of Go doesn't allow such mechanics, as Go was designed with go-routines in mind instead.

Solution A quick test with `cgo` yielded a viable solution to the problem: Using the C-routine `fork()` a fork was successful.

3.2 Shell instantiating and forwarding

Attempts in forwarding the client connection to a server-side shell's `stdin` and its `stdout` and `stderr` to the connection of the client were unsuccessful.

Solution One quick tests showed that hooking up the `std*` pipes to a local shell process with Go worked just fine. Therefore it was deemed feasible to transfer the entire interface to the client.

3.3 Participation of Mr. Schwarz

Up until this date, the participation of Mr Schwarz was remarkable little in terms of writing on the code base of the project. The present parties agreed on this matter.

Solution It was decided to give Mr Schwarz a choice of action: Either he starts to participate heavily in the project from now on or he opts out of the project entirely.

4 Old Business

- **Login attempts in Linux fail:** This problem was deemed lower priority, as Login works on the WSL and can still be dealt with in later stages of the project.

Next Meeting

Friday, 5th of April 2019, 13:00 in *ZL0.13, Lagerstrasse 45, Zürich*

Design and Implementation of an Alternative to SSH

Meeting Minutes

1 Attendees

Present: Stephan Neuhaus, Raphael Emberger, Kevin Schwarz

2 Initiation

The meeting took place on the *Friday, 12th of April 2019, 13:00* in *ZL0.13, Lagerstrasse 45, Zürich*. Raphael Emberger was responsible for the minutes.

3 Points of discussion

3.1 Participation of Mr. Schwarz

Mr Schwarz decided to opt out of the project because of time issues. Mr Neuhaus will therefore adapt the outline of the project.

3.2 Reading user data works

The new module to read user data via the `getpwnam(3)` API has been implemented using `cgo`. It can read all the required data (i.e. the user shell which wasn't supported in the go standard library).

3.3 Forking implemented, but causes problems

Forking has been implemented via `cgo` but after forking, the `net.Conn` object cannot be used by the child process. There is also the to further investigate, whether after forking a new process actually gets started, as a quick look at the processes didn't reveal that a fork has been processed.

Solution To counter this problem it is suggested to do the connection build up via `cgo` using the C-socket API. This returns an integer as a file descriptor, which shouldn't cause problems when forking.

3.4 Remote start and handling of a shell has issues

After successfully hooking up the channels from the client to the shell process, almost all mechanics work as expected with exception of missing characters like the `PS{1,2,3,4}` prompts.

Solution It is suggested to compare the environment variables of the child shell process and the usual terminals to see if there are deal breaking differences. Adjusting the child shells environment variables might fix the problem.

4 Old Business

- **Login attempts in Linux fail:** This problem was deemed lower priority, as Login works on the WSL and can still be dealt with in later stages of the project.

Next Meeting

Friday, 26th of April 2019, 13:00 in ZL0.13, Lagerstrasse 45, Zürich

Design and Implementation of an Alternative to SSH

Meeting Minutes

1 Attendees

Present: Stephan Neuhaus, Raphael Emberger

2 Initiation

The meeting took place on the *Friday, 26th of April 2019, 13:00* in *ZL0.13, Lagerstrasse 45, Zürich*. Raphael Emberger was responsible for the minutes.

3 Points of discussion

3.1 Login and shell usage

The remote login, starting and usage of a user shell works now. It still does not behave like intended, as there are warnings printed on the screen and every line written gets echoed back, but overall, it works.

3.2 Pty echoes stdin back to stdout on client

As described in the point above, when entering shell commands on the shell after remote login, the written lines get echoed back after hitting enter.

Solution The reason this was so, is because the terminal on the client was still in the *cooked* mode rather than the *raw* mode, which behaves differently from the default *cooked* mode, which reads line by line and catches and interprets signals like [Ctrl]+[C] or [Ctrl]+[D]. Setting the terminal into *raw* mode should solve the issue as explained in "The Linux Programming Interface".

3.3 Transfer of SIGWINCH/ioctl

As described in the first point of discussion, when dropping into the remote shell, there are warnings displayed about a problem with `ioctl`.

Solution In "The Linux Programming Interface" it is also mentioned that making the child the session leader would solve this issue.

3.4 Login with test user fails

Trying to login with the test user results in an error when starting the shell because of missing files.

Solution This issue was easily solved as the script for setting up the test user was faulty: It didn't properly create the home directory of the test user and since the process which dropped privilege after login didn't have root rights anymore, it couldn't enter the home directory. Therefore, the directory owner and permissions were amended.

3.5 Forking abandoned

3.4 of the last meeting suggested using the C-style sockets to pass the file descriptors to the child process, which should solve the issue with forking and still using the net.Conn object. This has been implemented and after some adjustment worked out well.

4 Old Business

- **Login attempts in Linux fail** This problem was deemed lower priority, as Login works on the WSL and can still be dealt with in later stages of the project.

Next Meeting

Friday, 3th of March 2019, 12:30 in ZL0.13, Lagerstrasse 45, Zürich

Design and Implementation of an Alternative to SSH

Meeting Minutes

1 Attendees

Present: Stephan Neuhaus, Raphael Emberger

2 Initiation

The meeting took place on the *Friday, 3th of May 2019, 13:00* in *ZL0.13, Lagerstrasse 45, Zürich*. Raphael Emberger was responsible for the minutes.

3 Points of discussion

3.1 Pty echoes stdin back to stdout on client

The suggested solution from last meetings 3.2 of setting the client terminal into raw mode worked out well and has since been implemented like this.

3.2 SockAddr

The x-package `unix` provides wrappers for common Unix API calls. Most of the handling with sockets could be transferred to their methods, except `getpeername(3)`, which was implemented in the old way using `CGo`. The reason being that the returned `SockAddr` interface from `accept(3)` and `getpeername(3)` couldn't be casted (or rather: type-asserted) to be used as `SockAddrInet4` struct.

Solution After some experiments and searching the internet, it was revealed that the type assertion had to use pointers instead.

3.3 Role of login in the application

Just relying on `login` to perform all the important steps to log a user in is not a viable solution, as this would make log in via keys impossible. Therefore, it was made so the server decides whether to attempt to log in via keys or just to call `login`. Key authentication has been implemented as well and doesn't rely on `login`.

3.4 Documentation and bug-fixes before new features

It was deemed better to stop implementing new features as time is about to run out. Therefore, it is suggested to invest 80% of the remaining time to write on the documentation and 20% on bug-fixing. Only after everything has been finished, can new features be implemented - if at all.

3.5 Documentation chapter renaming

The original template for the BSc thesis has chapters named "Theoretical Principles" and "Method" in them, which will be renamed to the more appropriate "Design" and "Implementation". The latter will not reflect the progress in a chronological order but instead order it point by point.

3.6 Login attempts in Linux fail

Relying on `login` finally solved the problem of logging in on a specific linux distribution. The log in via keys also works, as the problem was originally PEM dependent, which gets omitted when authenticating via keys.

Next Meeting

Friday, 10th of March 2019, 12:30 in ZL0.13, Lagerstrasse 45, Zürich

Design and Implementation of an Alternative to SSH

Meeting Minutes

1 Attendees

Present: Stephan Neuhaus, Raphael Emberger

2 Initiation

The meeting took place on the *Friday, 10th of May 2019, 12:20* in *ZL0.13, Lagerstrasse 45, Zürich*. Raphael Emberger was responsible for the minutes.

3 Points of discussion

3.1 File descriptors and Go

After resolving last week's issue with `SocketAddr(3.2)`, the same approach enabled us to get the fd of a connection: By type-asserting it as `*net.TCPConn`, which had a function that returns the underlying file, which in turn had a function for returning the underlying fd.

3.2 Privilege dropping

After authenticating the user via keys, the host process spawns a shell with the appropriate privileges. For more secure handling, it was deemed best to drop privilege for the entire host process after authenticating via keys. Despite running with root rights (to enable a user to login as any user), the process receives an exception when calling `unix.Setuid` or `unix.Setgid`. The reason was assumed to be related to Go's infamous Go-routines.

3.3 Current State of Documentation

The current state of the documentation is still insufficient. There are 4 weeks until the deadline. It is suggested to write more on the documentation.

Next Meeting

Friday, 17th of March 2019, 13:00 in *ZL0.13, Lagerstrasse 45, Zürich*

Design and Implementation of an Alternative to SSH

Meeting Minutes

1 Attendees

Present: Stephan Neuhaus, Raphael Emberger

2 Initiation

The meeting took place on the *Friday, 17th of May 2019, 12:30* in *ZL0.13, Lagerstrasse 45, Zürich*. Raphael Emberger was responsible for the minutes.

3 Points of discussion

3.1 Referencing Linux Man Pages

As the code base heavily relies on the Linux API, it was decided to include bibtex references to the manual pages to improve readability.

3.2 Clarification of Design of Public Key Cryptography

When describing public key cryptography, the current state of the documentation doesn't go into further detail other than describing the flow of actions when authenticating via public key cryptography. This part has to be improved.

3.3 Login as Root possible

The current state of the implementation allows a client to authenticate itself and log in as root on the server. This is a point that should be improved, but can be postponed for now.

Solution When using `login(1)` to authenticate and log in a user, the `-f` flag can be used to specify a specific user. The client could already ask the user before communication with the server, which user should be used when logging in. This is partially already implemented.

3.4 Keys remain in memory

The current state of the implementation of the client still retains the private key in memory when performing key authentication. This could be a possible vulnerability which could be used by a third party as an attack vector to obtain said key.

3.5 Keys not stored optimally

The public key up to today stored the full public key inside a directory structure in the root user's home directory. This could cause problems when storing keys and should be changed to the way `ssh` stores authorized public keys: By storing the authorized keys for a server-side user in hashed format in its home directory.

3.6 RSync not implemented yet

The layout of the tasks mentions `rsync` compatibility as both a required use-case for a passing grade and an optional extra feature. This has been amended by Mr Neuhaus by stating explicitly, that `rsync` compatibility is optional.

3.7 Separate bibliography for man page references

As references to the Linux manual pages take up a considerable part of the overall references, it was considered to split it up into a manual-only bibliography and a normal bibliography. This suggestion was rejected as it was deemed tolerable to have one single bibliography.

3.8 Picture for the Publication Tool

The official Publication Tool requires a picture to be attached to the abstract when handing it in to the online tool. As this thesis is centered around a CLI application, no picture has been made so far. It was suggested to simply take a screenshot of the application in operation.

Next Meeting

The next meeting doesn't have a set date, time or place, as it was deemed a better option to organize a new meeting whenever the need for one arises.

Design and Implementation of an Alternative to SSH

Meeting Minutes

1 Attendees

Present: Stephan Neuhaus, Raphael Emberger

2 Initiation

The meeting took place on the *Friday, 31th of May 2019, 12:30* in *ZL0.13, Lagerstrasse 45, Zürich*. Raphael Emberger was responsible for the minutes.

3 Points of discussion

3.1 Time Issues

Because of the lack of participation from a peer, Raphael was occupied with the labs of the AI2 course for over two days, which led to more pressure in finishing the thesis.

3.2 Project stopped working

When connecting to the server, the host application cannot open the connection from the file descriptor and throws a "Bad file-descriptor" error. This happens on the WSL.

Using a native Linux distribution solved this issue.

3.3 Formal Requirements

It was stressed to find out more about the formal requirements, lest the whole thesis might face tribulations due to formal details.

3.4 Attendance of Mr. Schwarz at the Presentation

It was suggested to include Mr. Schwarz in the presentation of the thesis as he was part of the team for a considerable amount of time. He has no obligation to present himself but may attend the presentation itself.

Next Meeting

There will not be any further meetings.

A.2. Others

Please refer to the USB-stick that has been handed in with this thesis. The Git repository for the code (private) is located at: <https://github.engineering.zhaw.ch/neut/oh-my-gosh> and the documentation (clear-net) at: <https://github.engineering.zhaw.ch/neut/oh-my-gosh-bericht>.