

Nondeterministic Asynchronous Dataflow in Isabelle/HOL

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Motivation

Context:

- Stream Processing: programs that compute (possibly) unbounded sequences of data (streams)
- A common problem in the industry
- Frameworks:
Apache Flink, Apache Samza, Apache Spark, Google Cloud Dataflow, and Timely Dataflow



- Why use frameworks?
 - Highly Parallel
 - Low latency (output as soon as possible)
 - Incremental computing (re-uses previous computations)

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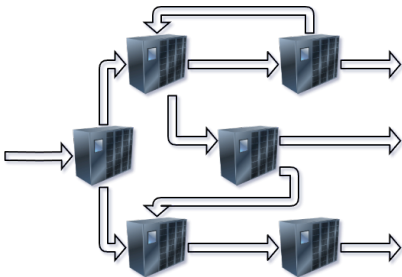
Our goal:

Mechanically Verify Timely Dataflow algorithms

A Good Foundation

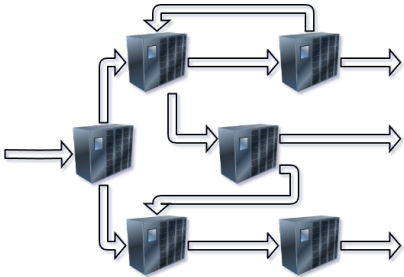
A Good Foundation

- Nondeterministic Asynchronous Dataflow

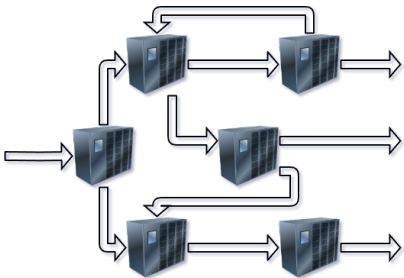


A Good Foundation

- Nondeterministic Asynchronous Dataflow
 - Dataflow: Directed graph of interconnected operators

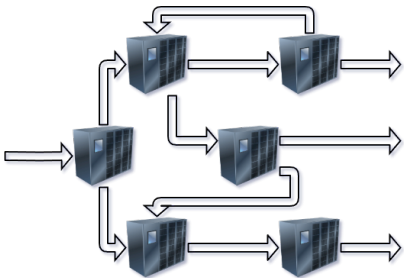


A Good Foundation



- Nondeterministic Asynchronous Dataflow
 - Dataflow: Directed graph of interconnected operators
 - Asynchronous:
 - Operators execute independently: processes without an orchestrator
 - Operators can freely communicate with the network (read/write); do silent computation steps
 - Networks are unbounded FIFO queues

A Good Foundation



- Nondeterministic Asynchronous Dataflow
 - Dataflow: Directed graph of interconnected operators
 - Asynchronous:
 - Operators execute independently: processes without an orchestrator
 - Operators can freely communicate with the network (read/write); do silent computation steps
 - Networks are unbounded FIFO queues
 - Nondeterministic:
 - Operators can make nondeterministic choices
 - Operators are relations between inputs and outputs sequences

The Algebra for Nondeterministic Asynchronous Dataflow

- Bergstra et al. presents an algebra for Nondeterministic Asynchronous Dataflow
- Primitives:
sequential and parallel composition; feedback loop...
- The 52 axioms
- An process calculus instance

Network Algebra for Asynchronous Dataflow*

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Isabelle/HOL Preliminaries

- Classical higher-order logic (HOL):
Simple Typed Lambda Calculus + axiom of choice + axiom of infinity + rank-1 polymorphism

Isabelle/HOL

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- Isabelle: A generic proof assistant



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Simple Typed Lambda Calculus + axiom of choice + axiom of infinity + rank-1 polymorphism
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- Isabelle/HOL: Isabelle's flavor of HOL

Why Isabelle/HOL?

- Codatatypes: (possibly) infinite data structures (e.g., lazy lists, streams)
- Corecursion: always eventually produces some codatatype constructor
- Coinductive predicate: infinite number of introduction rule applications
- Coinduction: reason about coinductive predicates

Operators as a Codatatype

Operators in Isabelle/HOL

```
codatatype (inputs: 'i, outputs: 'o, 'd) op =  
  Read 'i ('d  $\Rightarrow$  ('i, 'o, 'd) op) | Write (('i, 'o, 'd) op) 'o 'd  
  Silent ('i, 'o, 'd) op | Choice (('i, 'o, 'd) op) cset
```

- Type parameters: inputs/output ports; data
- Operator's actions
- Possibly infinite trees
- inputs/outputs: Sets of used ports

Examples 1

Uncommunicative operators

abbreviation

$\oslash \equiv \text{Choice } \{\}_c$

corec spin_op (\otimes) where

$\otimes = \text{Choice } ((\lambda_. \otimes)`_c \{()\}_c)$

corec silent_op (\odot) where

$\odot = \text{Silent } \odot$

lemma spin_op_code:

$\otimes = \text{Choice } \{\otimes\}_c$

- Quirk: any corecursive call guarded by the Choice must be applied using the proper map function
- They have the same meaning.

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Uncommunicative operators

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lemma spin_op_code:

$\otimes = \text{Choice } \{\otimes\}_c$

- Quirk: any corecursive call guarded by the Choice must be applied using the proper map function
- They have the same meaning. But are syntactically different ☹

Operators Equivalences: Motivation

An ideal equivalence relation

- Only equate operators that **behave** the same
- Useful reasoning principle

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An ideal equivalence relation

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- Milner's approach: Bisimilarity
- Based on labeled transition systems (LTS)
- Labels (actions): Read, Write, Silent (τ)

Intuition

Two operators are bisimilar if their corresponding transition systems can mutually simulate each other's transitions.

Label Transition System

datatype ('i,'o,'d) *IO* = Inp 'i 'd | Out 'o 'd | Tau

inductive step **where**

 step (Inp *p x*) (Read *p f*) (*f x*) | step (Out *q x*) (Write *op q x*) *op*
 | step Tau (Silent *op*) *op*
 | $op \in_c ops \implies \text{step } io \text{ } op \text{ } op' \implies \text{step } io \text{ } (\text{Choice } ops) \text{ } op'$

Operators Equivalences: Weak Bisimilarity

- foo

Asynchronous Dataflow Operators

- foo

Asynchronous Dataflow Properties

Basic network algebra properties

B1: $op_1 \parallel (op_2 \parallel op_3) \approx \text{map_op } \curvearrowright \curvearrowright (op_1 \parallel op_2) \parallel op_3$
 B2_1: $op \parallel (\mathcal{I} :: (0, 0, 'd) \text{ } op) \approx \text{map_op } \text{Inl } \text{Inl } op$
 B2_2: $(\mathcal{I} :: (0, 0, 'd) \text{ } op) \parallel op \approx \text{map_op } \text{Inr } \text{Inr } op$
 B3: $(op_1 \bullet op_2) \bullet op_3 \approx op_1 \bullet (op_2 \bullet op_3)$
 B4_1: $op \sqcap \bullet \mathcal{I} \approx op \sqcap$ B4_2: $\mathcal{I} \bullet \sqcap op \approx \sqcap op$
 B5: $(op_1 \parallel op_2) \bullet (op_3 \parallel op_4) \approx (op_1 \bullet op_3) \parallel (op_2 \bullet op_4)$
 B6: $\mathcal{I} \parallel \mathcal{I} \approx \mathcal{I}$ B7: $\mathcal{X} \bullet \mathcal{X} \approx \mathcal{I}$
 B8: $(\mathcal{X} :: ('i + 0, 0 + 'i, 'd) \text{ } op) \approx \text{map_op } \text{id } (\text{case_sum } \text{Inr } \text{Inl}) \mathcal{I}$
 B9: $\mathcal{X} \approx \text{map_op } \curvearrowright \curvearrowright (\mathcal{X} \parallel \mathcal{I}) \bullet \text{map_op } \text{id } \curvearrowright (\mathcal{I} \parallel \mathcal{X})$
 B10: $(\sqcap op_1 \parallel \sqcap op_2) \bullet \mathcal{X} \approx \mathcal{X} \bullet (op_2 \sqcap \parallel op_1 \sqcap)$
 F1: $\mathcal{I} \uparrow \approx (\mathcal{I} :: (0, 0, 'd) \text{ } op)$ F2: $\mathcal{X} \uparrow \approx \mathcal{I}$

R1: $\text{Inr } \vdash \text{inputs } op_1 \cap \text{defaults} = \{\} \implies \text{Inr } \vdash \text{outputs } op_1 \cap \text{defaults} = \{\} \implies$
 $op_2 \bullet (op_1 \uparrow) \approx ((op_2 \parallel \mathcal{I}) \bullet op_1) \uparrow$
 R2: $\text{Inr } \vdash \text{inputs } op_1 \cap \text{defaults} = \{\} \implies \text{Inr } \vdash \text{outputs } op_1 \cap \text{defaults} = \{\} \implies$
 $(op_1 \uparrow) \bullet op_2 \approx (op_1 \bullet (op_2 \parallel \mathcal{I})) \uparrow$
 R3: $\text{Inr } \vdash \text{inputs } op_2 \cap \text{defaults} = \{\} \implies \text{Inr } \vdash \text{outputs } op_2 \cap \text{defaults} = \{\} \implies$
 $op_1 \parallel (op_2 \uparrow) \approx (\text{map_op } \curvearrowright \curvearrowright (op_1 \parallel op_2)) \uparrow$
 R4: $\text{Inr } \vdash \text{inputs } op_1 \cap \text{defaults} = \{\} \implies \text{Inr } \vdash \text{outputs } op_1 \cap \text{defaults} = \{\} \implies$
 $\text{inputs } op_2 \cap \text{defaults} = \{\} \implies \text{outputs } op_2 \cap \text{defaults} = \{\} \implies$
 $(\sqcap op_1 \bullet (\mathcal{I} \parallel op_2)) \uparrow \approx ((\mathcal{I} \parallel op_2) \bullet op_1 \sqcap) \uparrow$
 R5: $\text{Inr } \vdash \text{inputs } op = \{\} \implies \text{Inr } \vdash \text{outputs } op = \{\} \implies$
 $\text{map_op } \text{Inl } \text{Inl } ((op :: ('i + 0, 'o + 0, 'd) \text{ } op) \uparrow) \approx op$
 R6: $\text{Inr } \vdash \text{inputs } op = \{\} \implies \text{Inr } \vdash \text{outputs } op = \{\} \implies$
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■ **Table 1** Basic network algebra properties

Basic network algebra properties

$$B1: op_1 \parallel (op_2 \parallel op_3) \approx \text{map_op } \hookrightarrow \hookrightarrow (op_1 \parallel op_2) \parallel op_3$$

$$B2_1: op \parallel (\mathcal{I} :: (0, 0, 'd) op) \approx \text{map_op } \text{Inl } \text{Inl } op$$

$$B2_2: (\mathcal{I} :: (0, 0, 'd) op) \parallel op \approx \text{map_op } \text{Inr } \text{Inr } op$$

$$B3: (op_1 \bullet op_2) \bullet op_3 \approx op_1 \bullet (op_2 \bullet op_3)$$

$$B4_1: op \sqcap \bullet \mathcal{I} \approx op \sqcap$$

$$B4_2: \mathcal{I} \bullet \sqcap op \approx \sqcap op$$

$$B5: (op_1 \parallel op_2) \bullet (op_3 \parallel op_4) \approx (op_1 \bullet op_3) \parallel (op_2 \bullet op_4)$$

$$B6: \mathcal{I} \parallel \mathcal{I} \approx \mathcal{I}$$

$$B7: \mathcal{X} \bullet \mathcal{X} \approx \mathcal{I}$$

$$B8: (\mathcal{X} :: ('i + 0, 0 + 'i, 'd) op) \approx \text{map_op } \text{id } (\text{case_sum } \text{Inr } \text{Inl}) \mathcal{I}$$

$$B9: \mathcal{X} \approx \text{map_op } \hookrightarrow \hookrightarrow (\mathcal{X} \parallel \mathcal{I}) \bullet \text{map_op } \text{id } \hookrightarrow (\mathcal{I} \parallel \mathcal{X})$$

$$B10: (\sqcap op_1 \parallel \sqcap op_2) \bullet \mathcal{X} \approx \mathcal{X} \bullet (op_2 \sqcap \parallel op_1 \sqcap)$$

$$F1: \mathcal{I} \uparrow \approx (\mathcal{I} :: (0, 0, 'd) op)$$

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$$R1: \text{Inr } \dot{\vdash} \text{ inputs } op_1 \cap \text{ defaults } = \{\} \implies \text{Inr } \dot{\vdash} \text{ outputs } op_1 \cap \text{ defaults } = \{\} \implies op_2 \bullet (op_1 \uparrow) \approx ((op_2 \parallel \mathcal{I}) \bullet op_1) \uparrow$$

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$$R4: \text{Inr } \dot{\vdash} \text{ inputs } op_1 \cap \text{ defaults } = \{\} \implies \text{Inr } \dot{\vdash} \text{ outputs } op_1 \cap \text{ defaults } = \{\} \implies \text{inputs } op_2 \cap \text{ defaults } = \{\} \implies \text{outputs } op_2 \cap \text{ defaults } = \{\} \implies (\sqcap op_1 \bullet (\mathcal{I} \parallel op_2)) \uparrow \approx ((\mathcal{I} \parallel op_2) \bullet op_1 \sqcap) \uparrow$$

$$R5: \text{Inr } \dot{\vdash} \text{ inputs } op = \{\} \implies \text{Inr } \dot{\vdash} \text{ outputs } op = \{\} \implies \text{map_op } \text{Inl } \text{Inl } ((op :: ('i + 0, 'o + 0, 'd) op) \uparrow) \approx op$$

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$$R5: \text{Inr } \curvearrowright \text{inputs } op = \{\} \implies \text{Inr } \curvearrowright \text{outputs } op = \{\} \implies \text{map_op } \text{Inl } \text{Inl } ((op :: ('i + 0, 'o + 0, 'd) op) \uparrow) \approx op$$

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$$\text{B2_2: } (\mathcal{I} :: (\theta, \theta, 'd) \text{ op}) \parallel \text{op} \approx \text{map_op lnr lnr op}$$
$$\text{B3: } (op_1 \bullet op_2) \bullet op_3 \approx op_1 \bullet (op_2 \bullet op_3)$$
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$$\text{F1: } \mathcal{I} \uparrow \approx (\mathcal{I} :: (\theta, \theta, 'd) \text{ op}) \qquad \text{F2: } \mathcal{X} \uparrow \approx \mathcal{I}$$
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$$\begin{aligned} \text{R4: } \text{Inr } \vdash \text{inputs } op_1 \cap \text{defaults} = \{\} &\implies \text{Inr } \vdash \text{outputs } op_1 \cap \text{defaults} = \{\} \implies \\ \text{inputs } op_2 \cap \text{defaults} = \{\} &\implies \text{outputs } op_2 \cap \text{defaults} = \{\} \implies \\ (\neg op_1 \bullet (\mathcal{I} \parallel op_2)) \uparrow \approx ((\mathcal{I} \parallel op_2) \bullet op_1) \uparrow & \end{aligned}$$
$$\text{R5: } \text{Inr } \dot{\vdash} \text{ inputs } op = \{\} \implies \text{Inr } \dot{\vdash} \text{ outputs } op = \{\} \implies \text{map_op } \text{Inl } \text{Inl } ((op :: ('i + \theta, 'o + \theta, 'd) \text{ op}) \uparrow) \approx op$$
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Conclusion

- Isabelle/HOL has a good tool set to formalize and reason about stream processing:
 - Codatatypes, coinductive predicates, corecursion with friends, reasoning up to friends (congruence),
 - Coinduction up to congruence principle is automatically derived for codatatypes (but not for coinductive principles)
- Next step: Feedback loop

Questions, comments and suggestions