

Nondeterministic Asynchronous Dataflow in Isabelle/HOL

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Motivation

Context:

- Stream Processing: programs that compute (possibly) unbounded sequences of data (streams)
- A common problem in the industry
- Frameworks:
Apache Flink, Apache Samza, Apache Spark, Google Cloud Dataflow, and Timely Dataflow



- Why use frameworks?
 - Highly Parallel
 - Low latency (output as soon as possible)
 - Incremental computing (re-uses previous computations)

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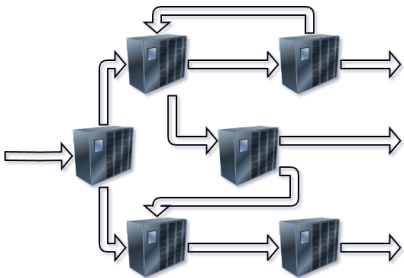
Our goal:

Mechanically Verify Timely Dataflow algorithms

A Good Foundation

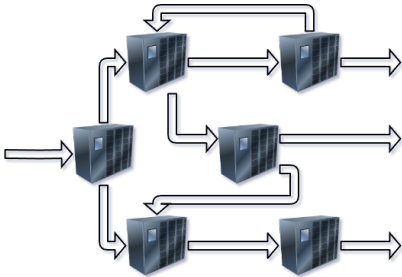
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- Nondeterministic Asynchronous Dataflow

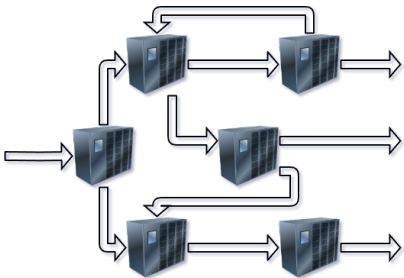


A Good Foundation

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 - Dataflow: Directed graph of interconnected operators

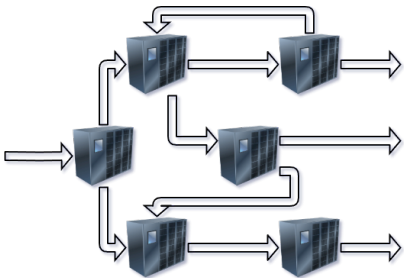


A Good Foundation



- Nondeterministic Asynchronous Dataflow
 - Dataflow: Directed graph of interconnected operators
 - Asynchronous:
 - Operators execute independently: processes without an orchestrator
 - Operators can freely communicate with the network (read/write); do silent computation steps
 - Networks are unbounded FIFO queues

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- Nondeterministic Asynchronous Dataflow
 - Dataflow: Directed graph of interconnected operators
 - Asynchronous:
 - Operators execute independently: processes without an orchestrator
 - Operators can freely communicate with the network (read/write); do silent computation steps
 - Networks are unbounded FIFO queues
 - Nondeterministic:
 - Operators can make nondeterministic choices
 - Operators are relations between inputs and outputs sequences

The Algebra for Nondeterministic Asynchronous Dataflow

- Bergstra et al. presents an algebra for Nondeterministic Asynchronous Dataflow
- Primitives:
sequential and parallel composition; feedback loop...
- The 52 axioms
- An process calculus instance

Network Algebra for Asynchronous Dataflow*

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Isabelle/HOL Preliminaries

- Classical higher-order logic (HOL):
Simple Typed Lambda Calculus + axiom of choice + axiom of infinity + rank-1 polymorphism

Isabelle/HOL

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- Isabelle/HOL: Isabelle's flavor of HOL

Why Isabelle/HOL?

- Codatatypes: (possibly) infinite data structures (e.g., lazy lists, streams)
- Corecursion: always eventually produces some codatatype constructor
- Coinductive predicate: infinite number of introduction rule applications
- Coinduction: reason about coinductive predicates

Operators as a Codatatype

Operators in Isabelle/HOL

```
codatatype (inputs: 'i, outputs: 'o, 'd) op =  
  Read 'i ('d  $\Rightarrow$  ('i, 'o, 'd) op) | Write (('i, 'o, 'd) op) 'o 'd  
  Silent ('i, 'o, 'd) op | Choice (('i, 'o, 'd) op) cset
```

- Type parameters: inputs/output ports; data
- Operator's actions
- Possibly infinite trees
- inputs/outputs: Sets of used ports

Examples 1

Uncommunicative operators

abbreviation

$$\odot \equiv \text{Choice } \{\}_c$$

corec spin_op (\otimes) where

$$\otimes = \text{Choice } ((\lambda_. \otimes)`_c \{()\}_c)$$

corec silent_op (\odot) where

$$\odot = \text{Silent } \odot$$

lemma spin_op_code:

$$\otimes = \text{Choice } \{\otimes\}_c$$

- Quirk: any corecursive call guarded by the Choice must be applied using the proper map function
- They have the same meaning.

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lemma spin_op_code:

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- Quirk: any corecursive call guarded by the Choice must be applied using the proper map function
- They have the same meaning. But are syntactically different ☹

Operators Equivalences: Motivation

An ideal equivalence relation

- Only equate operators that **behave** the same
- Useful reasoning principle

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An ideal equivalence relation

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- Milner's approach: Bisimilarity
- Based on labeled transition systems (LTS)
- Labels (actions): Read, Write, Silent (τ)

Intuition

Two operators are bisimilar if their corresponding transition systems can mutually simulate each other's transitions.

Label Transition System

datatype ('i,'o,'d) *IO* = Inp 'i 'd | Out 'o 'd | Tau

inductive step **where**

step (Inp *p x*) (Read *p f*) (*f x*) | step (Out *q x*) (Write *op q x*) *op*

| step Tau (Silent *op*) *op*

| $op \in_c ops \implies \text{step } io \text{ } op \text{ } op' \implies \text{step } io \text{ (Choice } ops) \text{ } op'$

Operators Equivalences: Strong Bisimilarity

Simulates

definition sim where

$$\text{sim } R \text{ } op_1 \text{ } op_2 = (\forall io \text{ } op'_1. \text{step } io \text{ } op_1 \text{ } op'_1 \longrightarrow (\exists op'_2. \text{step } io \text{ } op_2 \text{ } op'_2 \wedge R \text{ } op'_1 \text{ } op'_2))$$

Strong Bisimilarity

Operators Equivalences: Weak Bisimilarity

- foo

Asynchronous Dataflow Operators

- foo

Asynchronous Dataflow Properties

Basic network algebra properties

$$\text{B1: } op_1 \parallel (op_2 \parallel op_3) \approx \text{map_op} \curvearrowright \curvearrowright (op_1 \parallel op_2) \parallel op_3$$
$$\text{B2_1: } op \parallel (\mathcal{I} :: (\theta, \theta, 'd) \text{ } op) \approx \text{map_op } \text{Inl } \text{Inl } op$$
$$\text{B2_2: } (\mathcal{I} :: (\theta, \theta, 'd) \text{ op}) \parallel \text{op} \approx \text{map_op lnr lnr op}$$
$$\text{B3: } (op_1 \bullet op_2) \bullet op_3 \approx op_1 \bullet (op_2 \bullet op_3)$$
$$\text{B4_1: } op \vdash \bullet \mathcal{I} \approx op \vdash$$
$$\text{B4_2: } \mathcal{I} \bullet \dashv_{op} \approx \dashv_{op}$$
$$\text{B5: } (op_1 \parallel op_2) \bullet (op_3 \parallel op_4) \approx (op_1 \bullet op_3) \parallel (op_2 \bullet op_4)$$

B6: $\mathcal{I} \parallel \mathcal{I} \approx \mathcal{I}$

B7: $\mathcal{X} \bullet \mathcal{X} \approx \mathcal{I}$

$$\text{B8: } (\mathcal{X} :: ('i + \theta, \theta + 'i, 'd) \text{ op}) \approx \text{map_op id (case_sum lnr lnl)} \mathcal{I}$$
$$\text{B9: } \mathcal{X} \approx \text{map op } \hookrightarrow \hookrightarrow (\mathcal{X} \parallel \mathcal{I}) \bullet \text{map op id } \hookrightarrow (\mathcal{I} \parallel \mathcal{X})$$
$$\text{B10: } (\dashv op_1 \parallel \dashv op_2) \bullet \mathcal{X} \approx \mathcal{X} \bullet (op_2 \vdash \parallel op_1 \vdash)$$
$$\text{F1: } \mathcal{I} \uparrow \approx (\mathcal{I} :: (0, 0, 'd) \text{ op})$$
$$\text{F2: } \mathcal{X} \uparrow \approx \mathcal{I}$$
$$\text{R1: } \text{Inr } \dot{\vdash} \text{ inputs } op_1 \cap \text{ defaults} = \{\} \implies \text{Inr } \dot{\vdash} \text{ outputs } op_1 \cap \text{ defaults} = \{\} \implies op_2 \bullet (op_1 \uparrow) \approx ((op_2 \parallel \mathcal{I}) \bullet op_1) \uparrow$$
$$\text{R2: } \text{Inr } \dot{\vdash} \text{ inputs } op_1 \cap \text{defaults} = \{\} \implies \text{Inr } \dot{\vdash} \text{ outputs } op_1 \cap \text{defaults} = \{\} \implies (op_1 \uparrow) \bullet op_2 \approx (op_1 \bullet (op_2 \parallel \mathcal{I})) \uparrow$$
$$\text{R3: } \text{Inr} \vdash \text{inputs } op_2 \cap \text{defaults} = \{\} \implies \text{Inr} \vdash \text{outputs } op_2 \cap \text{defaults} = \{\} \implies op_1 \parallel (op_2 \uparrow) \approx (\text{map_op } \hookrightarrow \hookrightarrow (op_1 \parallel op_2)) \uparrow$$
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$$\text{R5: } \text{Inr } \dot{\vdash} \text{ inputs } op = \{\} \implies \text{Inr } \dot{\vdash} \text{ outputs } op = \{\} \implies \text{map_op } \text{Inl } \text{Inl } ((op :: ('i + \theta, 'o + \theta, 'd) \text{ op}) \uparrow) \approx op$$
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■ **Table 1** Basic network algebra properties

Basic network algebra properties

$$B1: op_1 \parallel (op_2 \parallel op_3) \approx \text{map_op } \curvearrowright \curvearrowright (op_1 \parallel op_2) \parallel op_3$$

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$$B3: (op_1 \bullet op_2) \bullet op_3 \approx op_1 \bullet (op_2 \bullet op_3)$$

$$B4_1: op \sqcap \bullet \mathcal{I} \approx op \sqcap$$

$$B4_2: \mathcal{I} \bullet \sqcap op \approx \sqcap op$$

$$B5: (op_1 \parallel op_2) \bullet (op_3 \parallel op_4) \approx (op_1 \bullet op_3) \parallel (op_2 \bullet op_4)$$

$$B6: \mathcal{I} \parallel \mathcal{I} \approx \mathcal{I}$$

$$B7: \mathcal{X} \bullet \mathcal{X} \approx \mathcal{I}$$

$$B8: (\mathcal{X} :: ('i + 0, 0 + 'i, 'd) op) \approx \text{map_op } \text{id } (\text{case_sum } \text{Inr } \text{Inl}) \mathcal{I}$$

$$B9: \mathcal{X} \approx \text{map_op } \curvearrowright \curvearrowright (\mathcal{X} \parallel \mathcal{I}) \bullet \text{map_op } \text{id } \curvearrowright (\mathcal{I} \parallel \mathcal{X})$$

$$B10: (\sqcap op_1 \parallel \sqcap op_2) \bullet \mathcal{X} \approx \mathcal{X} \bullet (op_2 \sqcap \parallel op_1 \sqcap)$$

$$F1: \mathcal{I} \uparrow \approx (\mathcal{I} :: (0, 0, 'd) op)$$

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$$R5: \text{Inr } \dot{\vdash} \text{ inputs } op = \{\} \implies \text{Inr } \dot{\vdash} \text{ outputs } op = \{\} \implies \text{map_op } \text{Inl } \text{Inl } ((op :: ('i + 0, 'o + 0, 'd) op) \uparrow) \approx op$$

$$R6: \text{Inr } \dot{\vdash} \text{ inputs } op = \{\} \implies \text{Inr } \dot{\vdash} \text{ outputs } op = \{\} \implies \text{Inr } \dot{\vdash} \text{Inl } \dot{\vdash} \text{ inputs } op = \{\} \implies \text{Inr } \dot{\vdash} \text{Inl } \dot{\vdash} \text{ outputs } op = \{\} \implies (op \uparrow) \uparrow \approx (\text{map_op } \curvearrowright \curvearrowright op) \uparrow$$

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$$\text{F1: } \mathcal{I} \uparrow \approx (\mathcal{I} :: (0, 0, 'd) \text{ op}) \qquad \text{F2: } \mathcal{X} \uparrow \approx \mathcal{I}$$
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Conclusion

- Isabelle/HOL has a good tool set to formalize and reason about stream processing:
 - Codatatypes, coinductive predicates, corecursion with friends, reasoning up to friends (congruence),
 - Coinduction up to congruence principle is automatically derived for codatatypes (but not for coinductive principles)
- Next step: Feedback loop

Questions, comments and suggestions