

**University of Stuttgart**

Institute of Parallel and  
Distributed Systems (IPVS)

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# **Net-based Applications: Introduction**

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(Based on the slides of Dr. Frank Dürr)

# Overview

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## Distributed Systems and Applications

- Definition
- Reasons for Distribution
- Logical Model
- Communication Models
- System Software
- Client/Server Model
- Application Architectures



# Definition (1)

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- *A distributed system is one I cannot get any work done because a machine I never heard of has crashed* – L. Lamport
- *A distributed system consists of a collection of autonomous computers linked by a computer network and equipped with distributed system software* – Coulouris, Dollimore, Kindberg
- *Distributed applications* are executed on distributed systems
  - Web-based applications
  - Booking and information systems
  - Enterprise resource planning (ERP) systems



# Definition (2)

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**Key concepts** of a distributed system are:

- Distribution transparency
  - **Access transparency:** Local and remote resources are accessed in the same way
  - **Location transparency:** The location of objects is not known to the users
  - **Replication transparency:** The number of copies that exist of an object is not known to the users
  - **Fragmentation transparency:** Objects are accessed without knowledge about any possible fragmentation
- Independent, autonomous machines



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# (Some) Reasons for Distribution

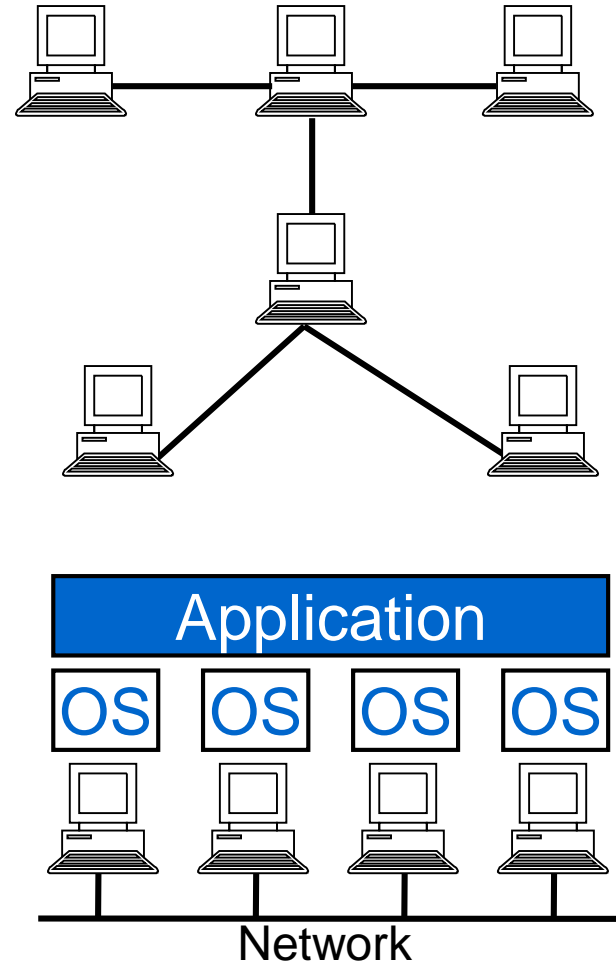
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- Resource sharing
  - Distributed printer / mainframe access
- Collaboration
  - Distributed file systems
  - Instant messenger
- Scalability
  - Centralized processing and storage does not scale well
  - Distributed processing and storage
    - Consistency and coordination problems
- Fault-Tolerance
  - Single computing / storage node results in single point of failure
  - Replication of data / pooled computing resources



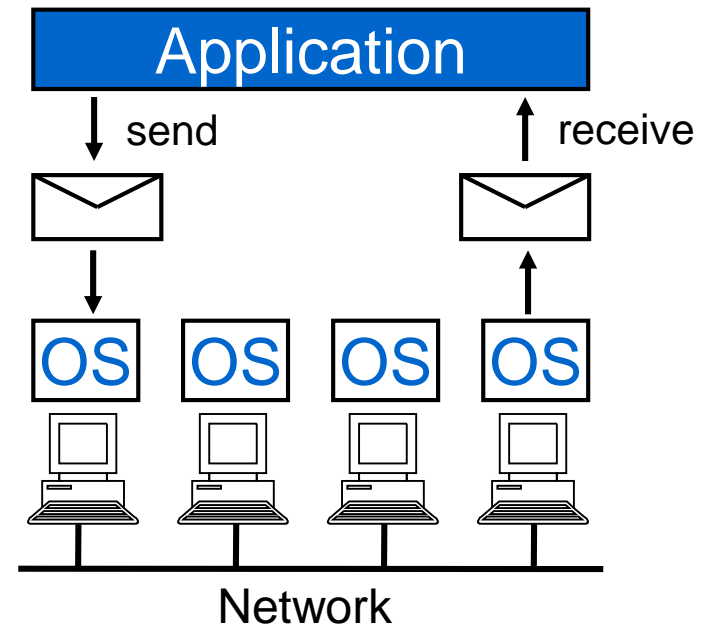
# Logical Model

- Distributed **system represented as graph** with nodes and edges
  - Nodes: autonomous computers
    - No shared physical memory
  - Edges: network connections
- **Simplified** logical model (no network and topological issues considered)
  - Nodes are connected by some network
  - Nodes may have different hardware and operating systems
  - Applications reside on nodes and **communicate across node boundaries**



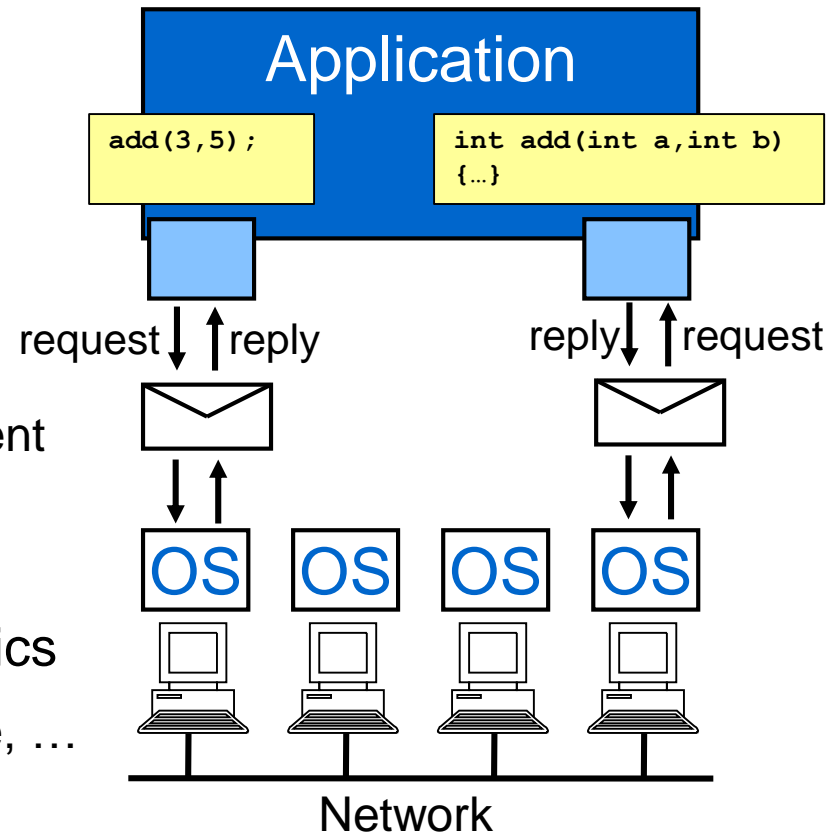
# Communication: Message Passing

- Data is exchanged explicitly
  - **send** and **receive** primitives
- Decoupling via **message buffers** at sender and / or receiver
  - Different **synchronization** points possible (blocking / unblocking)
- **Conversion of data** between heterogeneous end-systems necessary (marshalling)
- Programming abstraction differs from local interaction
  - Berkeley-Socket API
  - MPI



# Communication: Remote Procedure Call

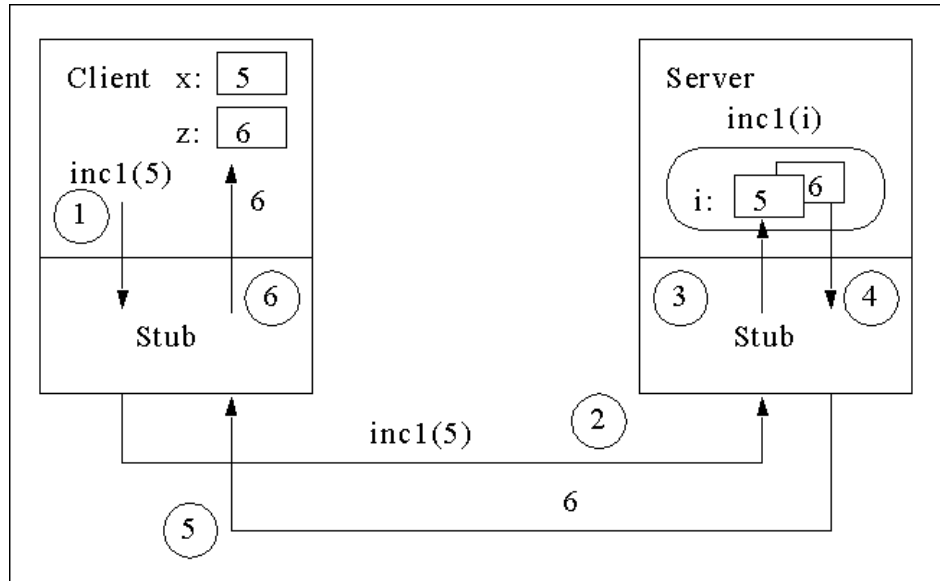
- Higher abstraction that **mimics local procedure call**
  - Similar to non-distributed interaction
  - Local proxies for remote entities
- **Automatic data-conversion**
  - Specification via machine independent language (Interface Definition Language, IDL)
- Different parameter passing semantics
  - By-value, by-reference, copy-restore, ...
- Examples
  - CORBA, DCOM, Java RMI





# Parameter Passing

## Call-by-Value



```
function inc1 (i : integer) : integer
```

```
begin
```

```
    i := i + 1;
```

```
    return (i);
```

```
end;
```

```
-- Call (client side):
```

```
    z := inc1(x);
```

Client stub : Copies the value of variable `x` as argument value into message

Server stub: Copies argument from message to variable `i`

Server stub: Copies value from variable `i` as result value into message

Client stub : Copies result from message to variable `z`



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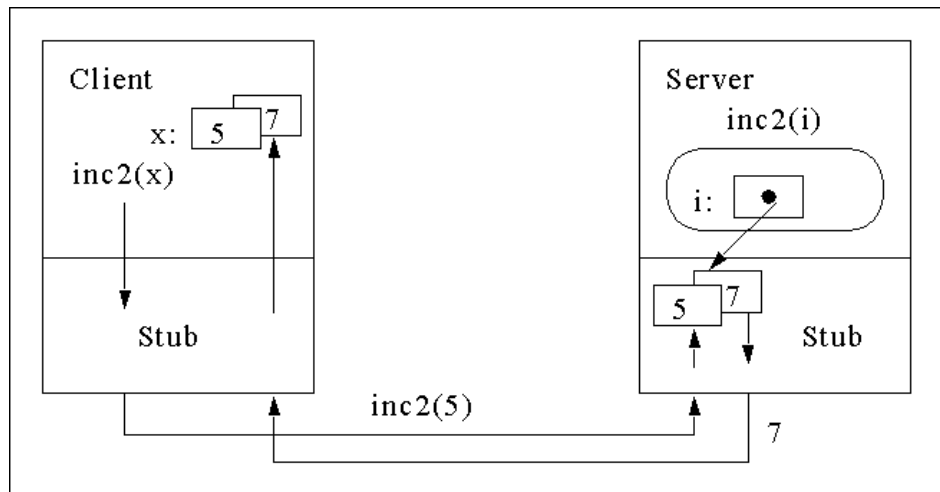
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# Parameter Passing

## Call-by-Reference

Possible implementation using a *call-by-copy/restore* approach:



```
function inc2 (var i : integer) begin  
    i := i + 2;  
end;  
  
-- Call (client side):  
inc2(x);
```

### Call:

Client stub : Copies variable `x` as argument into message

Server stub: Copies argument from message to stub variable

Passes *pointer to* stub variable to procedure

### Reply:

Server stub: Copies value from stub variable as result value into message

Client stub : Copies result from message to variable `x`



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# Difficulties with the Call-by-Copy/Restore Approach

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**program** test (output)

**var** a : integer;

**procedure** doubleinc

(var x, y : integer)

**begin**

    x := x + 1;

    y := y + 1;

**end;**

**begin**

    a := 0;

    doubleinc(a,a);

    writeln(a);

**end.**

**Problem:**

- Different results for local (a=2) and remote (a=1) execution
- Reason: Different copies for a single reference variable

**Alternative Implementation:** Server accesses parameter value remotely  $\Rightarrow$  huge overhead



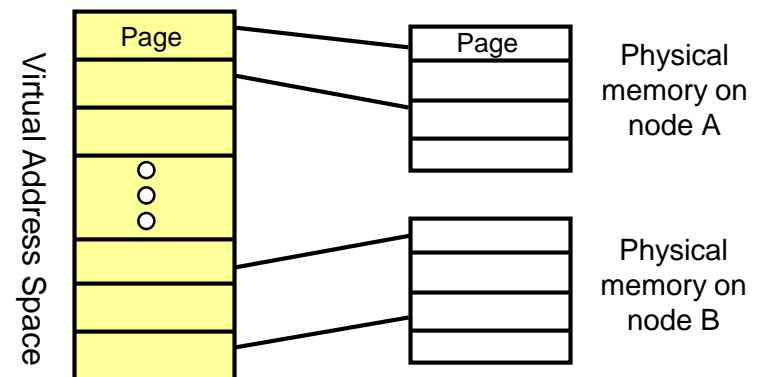
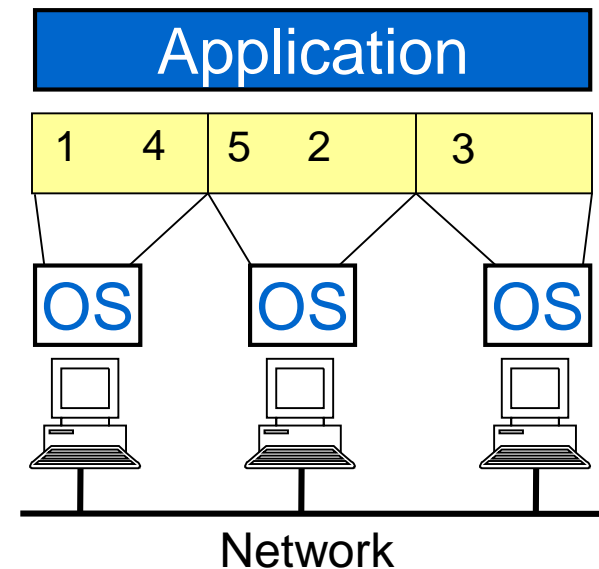
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# Communication: Distributed Shared Memory

- Distributed application shares a **common data space**
  - System realizes a global view on memory (granularity is page)
  - Communication via **read and write**
- Required memory pages are copied to local address space
  - Similar to demand-paging in OS
- Scalability and efficiency
  - Size bound by virtual address space size
  - Remote access requires copying of data → separation of execution and data-storage
  - Many optimizations possible (e.g. replicated read-only pages)



# System Software

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Many possible solutions to provide system software support for distributed applications

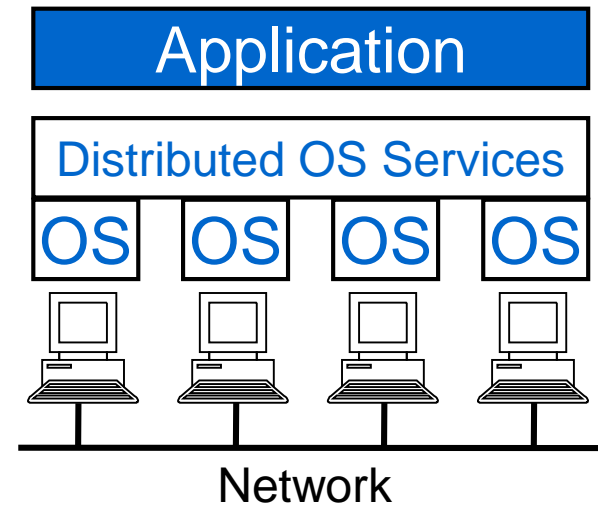
- **Distributed operating system**
  - Amoeba, Clouds
  - RPC, DSM
- **Network operating system** (OS with network extensions)
  - Unix, Windows
  - Sockets
- **Middleware** (on top of network operating system)
  - CORBA, DCOM, Java RMI, DCE
  - Remote procedure / method call
  - Distributed transactions



# Distributed Operating System

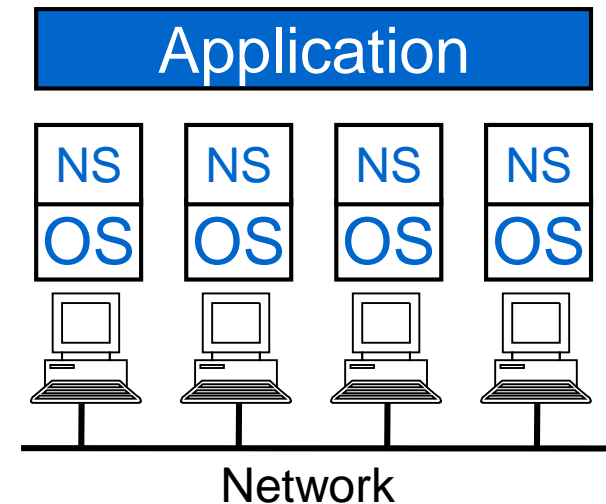
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- OS provides distributed virtual machine
  - Single system image
- Local kernel for local resource management enhanced by distributed operating system services
  - Distributed shared memory,
  - Task assignment (to machines),
  - Transparent file access,
  - Failure masking, ...
- Programming abstractions
  - High-level primitives as system calls
  - Used by applications and OS



# Network Operating System

- Operating system with network extensions
- Local kernel for local resource management with **basic network services**
  - Message exchange, e.g. sockets
- Additional services
  - User process
    - Remote procedure calls
  - Kernel extensions
    - Distributed file systems
- Programming abstractions
  - **Basic primitives for message exchange**
  - Cumbersome and error-prone



NS: Network Services

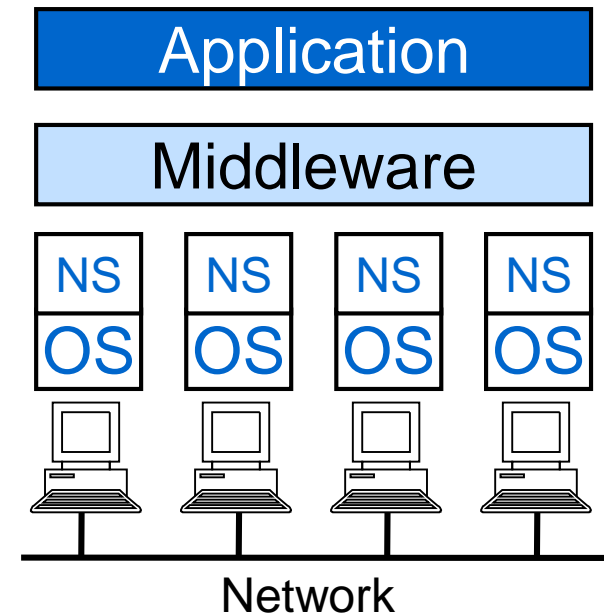


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# Communication Middleware

- Software layer **between OS and application**
  - Based on OS with network extension
  - Uses message exchange, e.g. sockets
- Objectives
  - **Uniform programming abstraction**
  - Mask heterogeneity
  - **Interoperability (common protocol)**
  - (Language independency)
- Programming abstractions
  - RPC, RMI
  - Message-oriented middleware (MOM)



NS: Network Services

## ➤ Brief RPC example

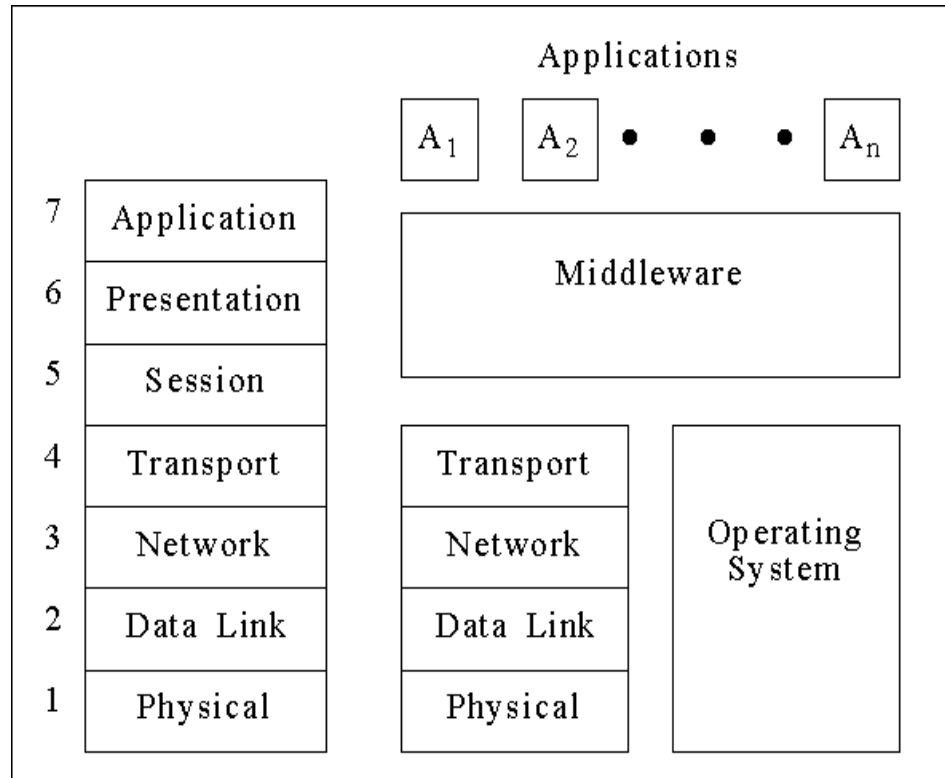


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# The Typical Architecture of a Distributed System with Middleware



The typical architecture of a distributed system in terms of the seven-layer OSI model.

## Transport System:

- Consists of layers 1 through 4
- Message-based inter-process communication
- No location transparency

## Middleware:

- Infrastructure for distributed applications
- Distribution transparency
- Examples:  
*Communication services, security services, directory services, time services, file services, transaction services, etc.*



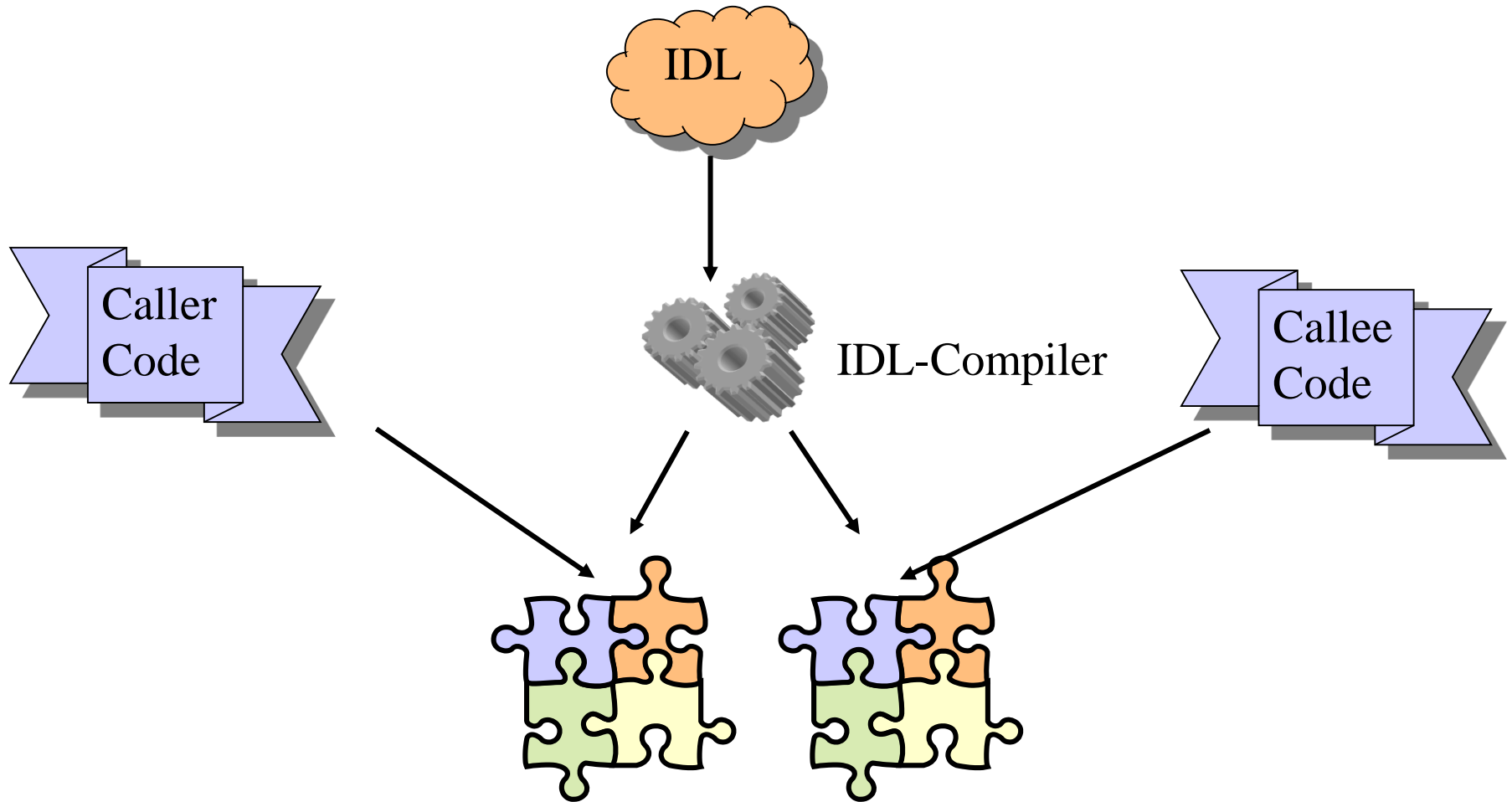
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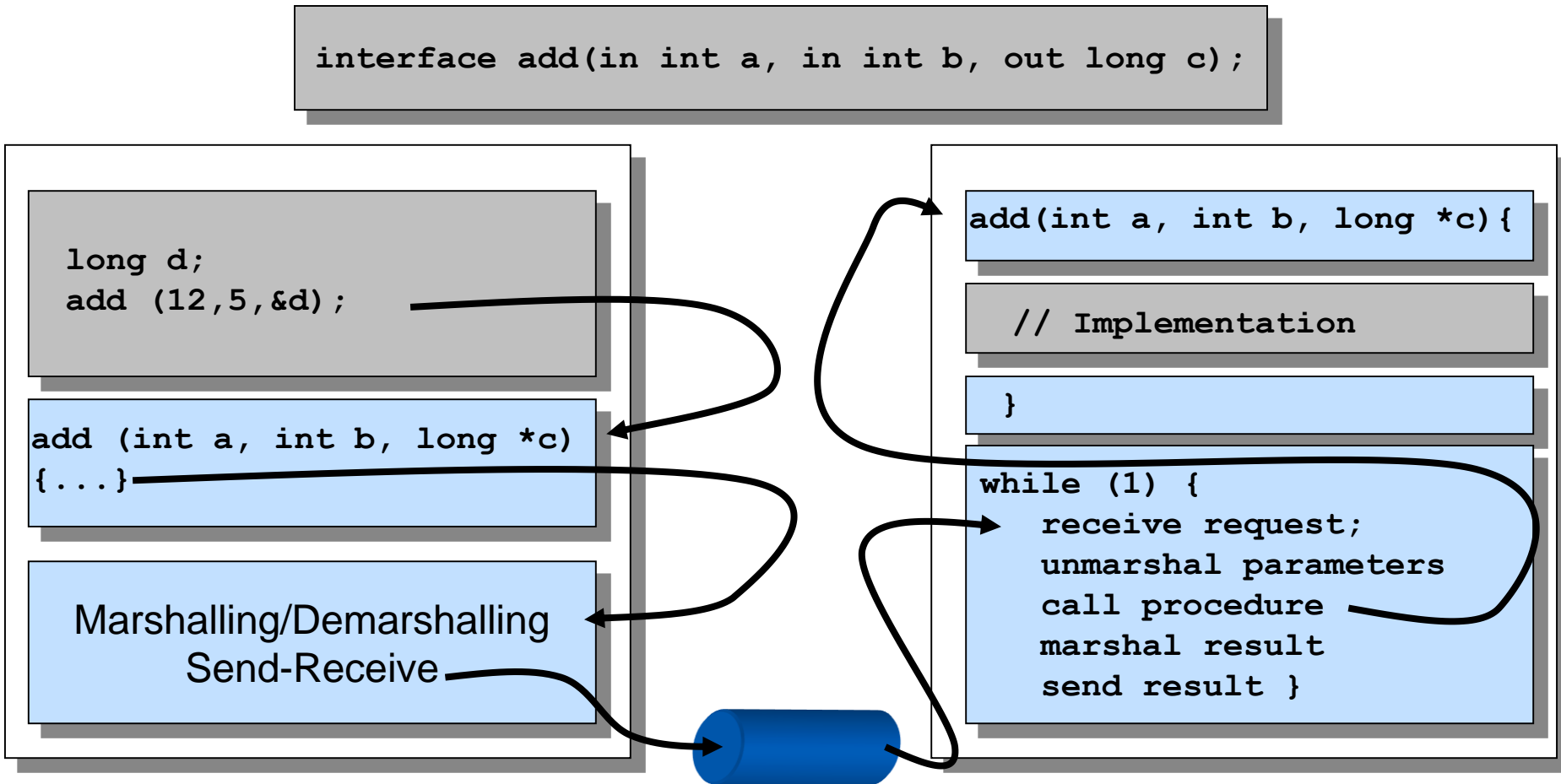
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# Example: Remote Procedure Call

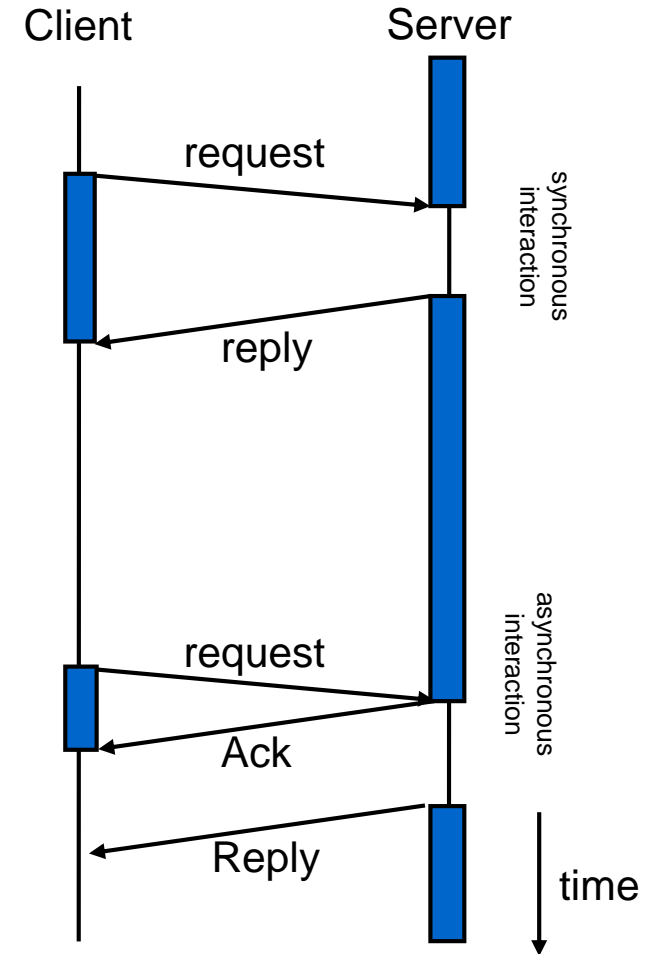


# Example: Remote Procedure Call



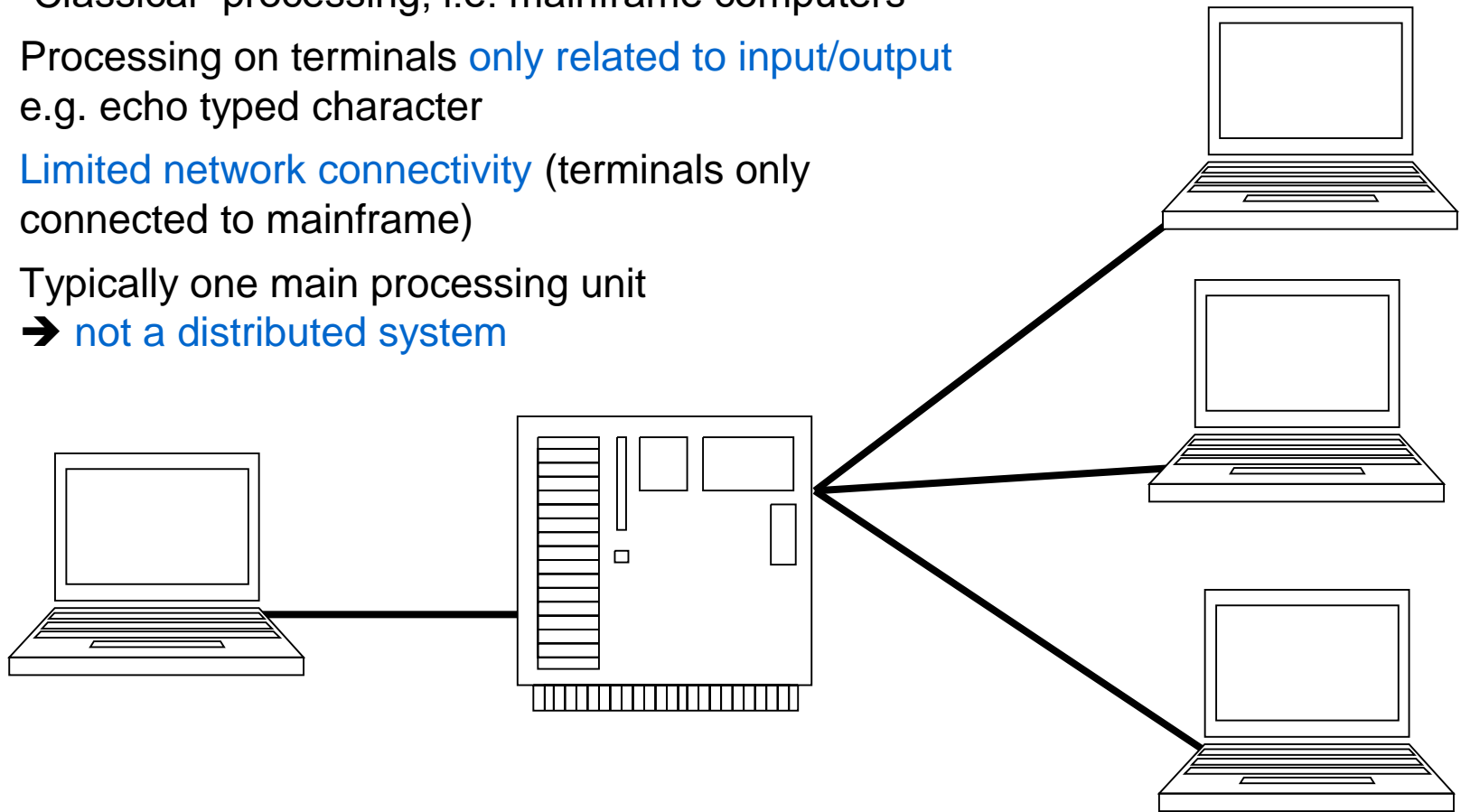
# Structuring Applications: Client-Server-Model

- Distributed applications can be structured in **two roles**
  - **Server**
    - **passive entity** (awaits incoming requests)
    - processes requests and returns results
  - **Client**
    - **active entity**
    - contacts server to request a service
- Different interaction models
  - Synchronous
  - Asynchronous
- Clear distinction between client and server in practice sometimes hard
  - Entities might **change their role**
  - There may be **thousands of clients and servers** (e.g. distributed objects)



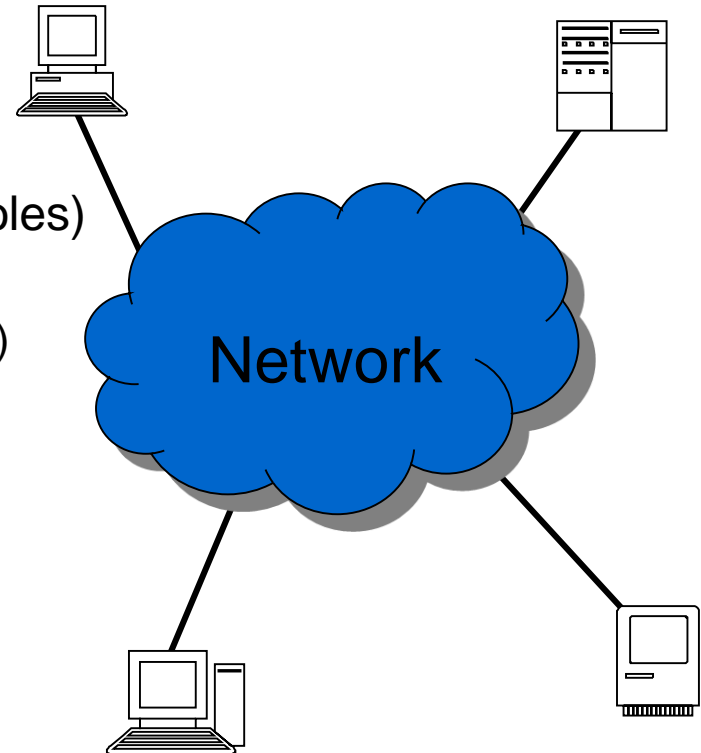
# Centralized Application Architecture

- “Classical” processing, i.e. mainframe computers
- Processing on terminals **only related to input/output**  
e.g. echo typed character
- **Limited network connectivity** (terminals only connected to mainframe)
- Typically one main processing unit  
→ **not a distributed system**



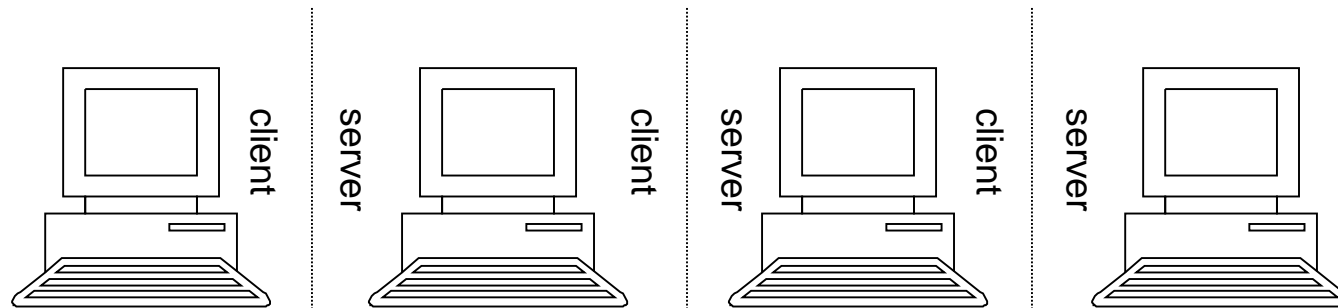
# Client/Server Application Architecture

- Heterogeneous system
  - Nodes: hardware, OS, language, ...
  - Edges: protocols, physical link, ...
- **Many clients and servers** (with changing roles)
  - No main processing system
  - Nodes may have equal rights (peer-to-peer)
- Possible (design) problems
  - Coordination hard to achieve
  - Performance bottlenecks
  - Single points of failure
  - No clear separation of logical tasks (e.g. presentation, processing, storage)



# N-Tier Application Architecture

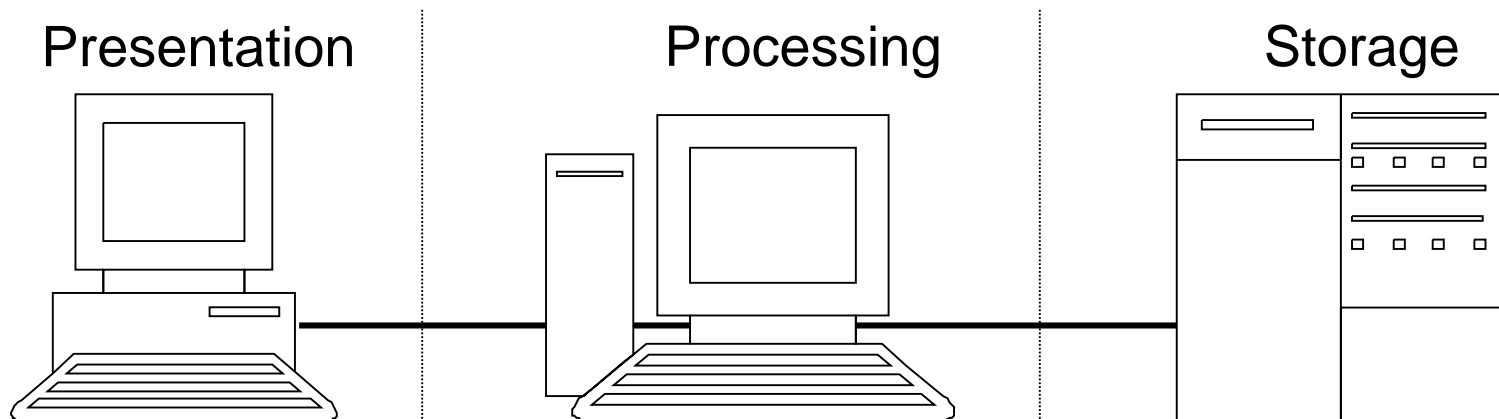
- Logical (and physical) **application structuring in N tiers** (aka multi-tier)
- Tiers provide a **distinct functionality** (e.g. store data)
- **Tier M is server for tier M-1 and client of tier M+1**
- Interaction only between “neighboring” tiers
  - Replacement of tiers less complex
  - Different implementations for individual tiers possible
- Design decision: How many tiers are needed? What is a good separation?



# 3-Tier Application Architecture

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- Widely used logical and physical separation
  - Tiers: presentation, processing, storage
- Still heterogeneous (e.g. multiple presentation tiers)
- Often centralization with respect to processing and storage
- How to distribute functionality among nodes in a distributed system?  
→ performance and fault-tolerance issues





# Presentation Tier

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- Responsibility
  - User interface to clients
  - Text based: shells, host masks
  - Graphical user interface (GUI)
    - Mouse and keyboard interaction
    - Windows, buttons, scrollbars, etc.
- Processing
  - Only related with presentation
    - Echoing keystrokes, requesting data on push button activity
- No data storage
- Clear distinction from application logic and data storage



# Logic (Application) Tier

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- Responsibility
  - Manipulation of data
  - Executing “business logic”
  - Access of persistent data from the database tier
  - Provides data to presentation tier
- No presentation of data
- Data storage
  - No persistent storage of data
  - Intermediate state may be captured for further processing



# Database Tier

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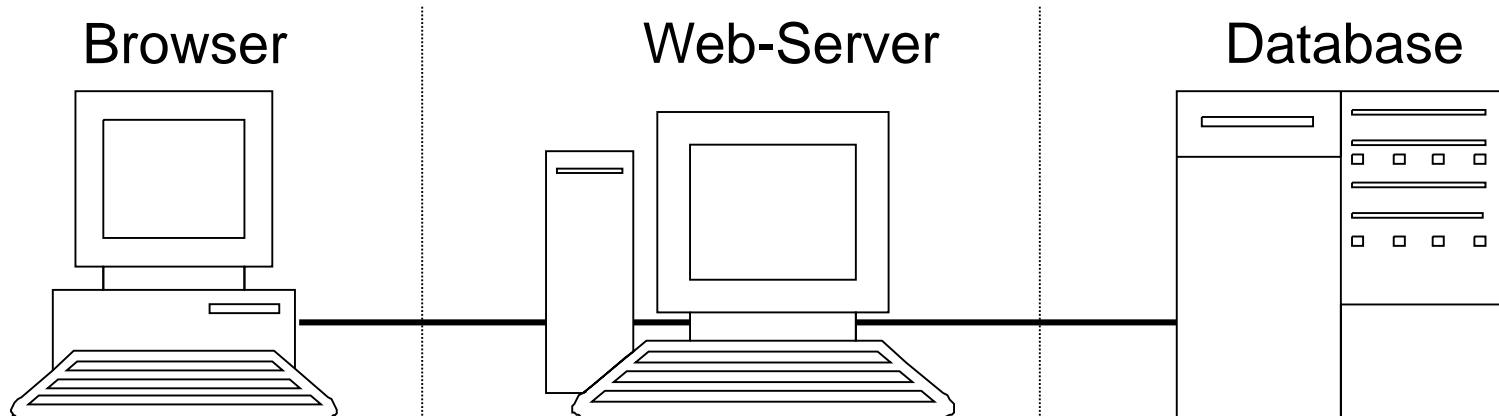
- Responsibility
  - Persistent storage of data
  - Concurrency control, i.e. transaction processing
  - Provides an appropriate interface to application tier
- No presentation
- Processing
  - Queries are processed and executed
  - No workflow or business logic



# Web-based Application Architecture

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- Special case of N-tier architecture
- Based on Internet technology, i.e. TCP/IP, HTTP
- Standard Components (browser, web server, DBMS)
- Customization: applets, servlets, database tables
- Suitable for Internet and intranets



# Net-based Applications ...

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- In this class we present and discuss important aspects of net-based applications:
  - Basic network programming: sockets, client/server ...
  - Web technologies: HTTP, JSP, Java Servlets, ...
  - XML Technology: specification, parsing, processing, ...
  - Basics of security: cryptography, certificates, ...

