CACHING

Write Policy

Write back: hit, standard, update mem on eviction

Through: update mem with cache

Write-allocate: miss, load and then write

No-write allocate: directly to memory

Placement

Direct: blocks map to one line, causes conflict misses

Fully-associative: blocks map anywhere, best hit rate, most complex

Set: a block can map to any line in a specific set

Lookup

Valid bit: being used, if 0, no assumptions!

Dirty (write back): modified

Mem addr: tag | index | offset

Index: set selection

Tag: line matching

Offset: within the block, pay attention to offset bit!

Replacement Policy

Least recently used: evict oldest used line in full set, recently used means use again soon

General Model

$$(s, E, b)$$

$\boldsymbol{s}$**:** # of set bits

$\boldsymbol{b}$**:** # of offset bits

$\boldsymbol{rest}, \boldsymbol{t}$**:** # of tag bits

$$S = 2^s, B = 2^b$$

$\boldsymbol{S}$**:** # of sets

$\boldsymbol{E}$**:** associativity, lines/set

$\boldsymbol{B}$**:** block size in bytes

Memory Hierarchy

Register, LX Cache SRAM, DRAM, SSD local net, HDD, tape Web

PROCESSES

Program: static instructions

Process: executing program

Job: used by shell process to manage children

Process groups: children under parent, direct shell children get own pgid

Orphan:

Exceptions

Event triggered

Sync: caused by program

Async: external, hardware, signals, system, causes context switch

Handlers: processes them

Context Switching

Caused by exceptions, managed by kernel, another process running

Process Model

Fork: creates new process exactly like parent

Exec: child switches state new process

Tree: remember shell exercises that you got all wrong

Signals

Async inter-process comms triggered by events

Process decides may decide response

Signal blocking: masks, pending signals… shell exercise print order question

Zombies

Process that terminated but still using resources

Reaping: using $wait$ and handle termination status

Concurrency

Inter-process, w/ sig handlers

Threads

Concurrent, share data, not regs and stack

Race conditions!

VIRTUAL MEMORY

Virtual maps to physical, one to many

Translated by memory management unit

Organized in $pages$, like cache blocks

“Unlimited” virtual memory

Make’s compiler job easier, allows for reuse of programs