

Assignments 3

Exercise 4

Given the function $F : (State \rightarrow State_{\perp}) \rightarrow (State \rightarrow State_{\perp})$ defined as follows:

$$F(f) = \text{cond}(\mathcal{B}[n > 0], f \circ \mathcal{S}[x := 2 * x; n := n - 1], id)$$

- (a) Give an explicit definition for $F(\lambda\sigma. \perp)$, $F^2(\lambda\sigma. \perp)$ and $F^3(\lambda\sigma. \perp)$.
- (b) From the results above, conjecture a general definition for $F^i(\lambda\sigma. \perp)$ where $i \geq 1$.
[Optional] Prove by induction on i that your conjecture is correct.
- (c) Give an explicit definition for $\bigsqcup_i F^i(\lambda\sigma. \perp)$.
- (d) Which is the least fixed point of F ? Justify your answer.
- (e) Given the above, compute the state resulting from the execution of the following program

$$x := 1; \text{while } n > 0 \text{ do } (x := 2 * x; n := n - 1)$$

under the initial stage $\sigma = [n \rightarrow 4]$.

a) An explicit definition for $F(\lambda\sigma. \perp)$, $F^2(\lambda\sigma. \perp)$ and $F^3(\lambda\sigma. \perp)$

Let $F : (State \rightarrow State_{\perp}) \rightarrow (State \rightarrow State_{\perp})$ be a function defined as follow

$$F(f) = \text{cond}(\mathcal{B}[n > 0], f \circ \mathcal{S}[S], id)$$

where $S = (x := 2 * x; n := n - 1)$.

Let us consider $f_0 : State \rightarrow State_{\perp}$, $f_0 = \lambda\sigma. \perp$ for all σ . We will define explicitly $f_1 = F(f_0)$, $f_1 = F^2(f_0)$ and $f_3 = F^3(f_0)$:

Definition of $F(f_0)$

Let us consider a function $f_1 : State \rightarrow State_{\perp}$ defined as follow

$$f_1 := F(f_0) = \text{cond}(\mathcal{B}[n > 0], f_0 \circ \mathcal{S}[S], id)$$

$$= \lambda\sigma. \begin{cases} f_0 \circ \mathcal{S}[S] \sigma & \mathcal{B}[n > 0] \sigma \\ id \sigma & \mathcal{B}[n \leq 0] \sigma \end{cases}$$

computation

$$= \lambda\sigma. \begin{cases} f_0 \circ \sigma[x \rightarrow 2 \cdot \sigma(x), n \rightarrow \sigma(n) - 1] & \sigma(n) > 0 \\ \sigma & \sigma(n) \leq 0 \end{cases}$$

↑
apply state σ
is not the same as composition.

but since $f_0 = \lambda\sigma. \perp \forall \sigma$, then $f_0 \circ \sigma[x \rightarrow 2 \cdot \sigma(x), n \rightarrow \sigma(n) - 1] = \perp$ for all σ . Therefore

$$f_1 = F(f_0) = \lambda\sigma. \begin{cases} \perp & \sigma(n) > 0 \\ \sigma & \sigma(n) \leq 0 \end{cases} \quad \checkmark \quad (1)$$

Definition of $F^2(f_0)$

Let us consider a function $f_2 : State \rightarrow State_\perp$ defined as follow

$$\begin{aligned} f_2 &:= F^2(f_0) = F(f_1) = \text{cond}(\mathcal{B}[\![n > 0]\!], f_1 \circ \mathcal{S}[\![S]\!], id) \\ &= \lambda\sigma. \begin{cases} f_1 \circ \mathcal{S}[\![S]\!] \sigma & \mathcal{B}[\![n > 0]\!]\sigma \\ id \sigma & \mathcal{B}[\![n \leq 0]\!]\sigma \end{cases} \\ &= \lambda\sigma. \begin{cases} f_1 \circ \sigma[x \rightarrow 2 \cdot \sigma(x), n \rightarrow \sigma(n) - 1] & \sigma(n) > 0 \\ \sigma & \sigma(n) \leq 0 \end{cases} \end{aligned}$$

we know now that (1) $f_1 = \lambda\sigma_1. \perp$ if $\sigma_1(n) > 0$ and $f_1 = \lambda\sigma_1. \sigma_1$ (the identity) if $\sigma_1(n) \leq 0$, being more precise $f_1 = \lambda\sigma_1. \sigma_1$ either $\sigma_1(n) = 0$ or $\sigma_1(n) < 0$, i.e.:

$$f_1 = \lambda\sigma_1. \begin{cases} \perp & \sigma_1(n) > 0 \\ \sigma_1 & \sigma_1(n) = 0 \\ \sigma_1 & \sigma_1(n) < 0 \end{cases}$$

So, in order to compose f_1 with $\sigma[x \rightarrow 2 \cdot \sigma(x), n \rightarrow \sigma(n) - 1]$, we have to distinguish two cases when $\sigma_1(n) > 0$ and $\sigma_1(n) = 0$, i.e. when $\sigma(n) - 1 > 0$, and $\sigma(n) - 1 = 0$, respectively.

The first one is actually easy, $\sigma_1(n) > 0 \iff \sigma(n) - 1 > 0 \iff \sigma(n) > 1$ then

$$f_2 = \lambda\sigma. \perp, \quad \sigma(n) > 1$$

In the second one $\sigma_1(n) = 0 \iff \sigma(n) - 1 = 0 \iff \sigma(n) = 1$, then

$$f_2 = \lambda\sigma. \sigma[x \rightarrow 2 \cdot \sigma(x), n \rightarrow 0], \quad \sigma(n) = 1$$

And finally in the case $\sigma(n) \leq 0$, $f_2 = \lambda\sigma. \sigma$. Therefore

$$f_2 = \lambda\sigma. \begin{cases} \perp & \sigma(n) > 1 \\ \sigma[x \rightarrow 2 \cdot \sigma(x), n \rightarrow 0] & \sigma(n) = 1 \\ \sigma & \sigma(n) \leq 0 \end{cases} \quad \checkmark \quad (2)$$

Definition of $F^3(f_0)$

Let us consider a function $f_3 : State \longrightarrow State_\perp$ defined as follow

$$\begin{aligned}
 f_3 &:= F^3(f_0) = F^2(f_1) = F(f_2) = \text{cond}(\mathcal{B}[\![n > 0]\!], f_2 \circ \mathcal{S}[\![S]\!], id) \\
 &= \lambda\sigma. \begin{cases} f_2 \circ \mathcal{S}[\![S]\!] \sigma & \mathcal{B}[\![n > 0]\!]\sigma \\ id \sigma & \mathcal{B}[\![n \leq 0]\!]\sigma \end{cases} \\
 &= \lambda\sigma. \begin{cases} f_2 \circ \sigma[x \rightarrow 2 \cdot \sigma(x), n \rightarrow \sigma(n) - 1] & \sigma(n) > 0 \\ \sigma & \sigma(n) \leq 0 \end{cases}
 \end{aligned}$$

In order to proceed with the definition of f_3 we will think pretty similar to the way when we defined f_2 , we know that (2) $f_2 = \lambda\sigma_2. \perp$ if $\sigma_2(n) > 1$, $f_2 = \lambda\sigma_2. \sigma_2[x \rightarrow 2 \cdot \sigma_2(x), n \rightarrow 0]$ if $\sigma_2(n) = 1$ and $f_2 = \lambda\sigma_2. \sigma_2$ (the identity) if $\sigma_2(n) \leq 0$, being more precise $f_2 = \lambda\sigma_2. \sigma_2$ either $\sigma_2(n) = 0$ or $\sigma_2(n) < 0$, i.e.:

$$f_2 = \lambda\sigma_2. \begin{cases} \perp & \sigma_2(n) > 1 \\ \sigma_2[x \rightarrow 2 \cdot \sigma_2(x), n \rightarrow 0] & \sigma_2(n) = 1 \\ \sigma_2 & \sigma_2(n) = 0 \\ \sigma_2 & \sigma_2(n) < 0 \end{cases}$$

Again, when $\sigma(n) > 0$, we will compose f_2 with $\sigma[x \rightarrow 2 \cdot \sigma(x), n \rightarrow \sigma(n) - 1]$. Then, we have to distinguish three cases when $\sigma_2(n) > 1$, $\sigma_2(n) = 1$ and $\sigma_2(n) = 0$, i.e. when $\sigma(n) - 1 > 1$, $\sigma(n) - 1 = 1$ and $\sigma(n) - 1 = 0$, respectively.

The first one is actually easy, $\sigma_2(n) > 1 \iff \sigma(n) - 1 > 1 \iff \sigma(n) > 2$ then

$$f_3 = \lambda\sigma. \perp, \quad \sigma(n) > 2$$

The third one $\sigma_2(n) = 0 \iff \sigma(n) - 1 = 0 \iff \sigma(n) = 1$ then

$$f_3 = \lambda\sigma. \sigma[x \rightarrow 2 \cdot \sigma(x), n \rightarrow 0], \quad \sigma(n) = 1$$

But in the case when $\sigma_2(n) = 1$, we note that in addition to map n to 2 by σ , because $\sigma_2(n) = 1 \iff \sigma(n) - 1 = 1 \iff \sigma(n) = 2$, we have to compose

$$(\lambda\sigma_2. \sigma_2[x \rightarrow 2 \cdot \sigma_2(x), n \rightarrow 0]) \circ (\lambda\sigma. \sigma[x \rightarrow 2 \cdot \sigma(x), n \rightarrow 2])$$

In other words,

$$\begin{aligned}
& (\lambda\sigma_2.\sigma_2[x \rightarrow 2 \cdot \sigma_2(x), n \rightarrow 0]) \circ (\lambda\sigma.\sigma[x \rightarrow 2 \cdot \sigma(x), n \rightarrow 2]) \\
&= \lambda\sigma.\sigma[x \rightarrow 2 \cdot \sigma(x), n \rightarrow 0][x \rightarrow 2 \cdot \sigma(x), n \rightarrow 2] \\
&= \lambda\sigma.\sigma[x \rightarrow 2 \cdot (2 \cdot \sigma(x)), n \rightarrow 0] \\
&= \lambda\sigma.\sigma[x \rightarrow 2 \cdot 2 \cdot \sigma(x), n \rightarrow 0] \\
&= \lambda\sigma.\sigma[x \rightarrow 2^2 \cdot \sigma(x), n \rightarrow 0]
\end{aligned}$$

And finally, in the case $\sigma(n) \leq 0$, $f_2 = \lambda\sigma.\sigma$. Therefore

$$f_3 = \lambda\sigma. \begin{cases} \perp & \sigma(n) > 2 \\ \sigma[x \rightarrow 2^2 \cdot \sigma(x), n \rightarrow 0] & \sigma(n) = 2 \\ \sigma[x \rightarrow 2 \cdot \sigma(x), n \rightarrow 0] & \sigma(n) = 1 \\ \sigma & \sigma(n) \leq 0 \end{cases} \quad (3)$$

b) An explicit definition for $F^i(\lambda\sigma. \perp)$.

Let us consider a function $f_i : State \rightarrow State_\perp$ defined as follow

$$f_i := F^i(\lambda\sigma. \perp) = \lambda\sigma. \begin{cases} \perp & \sigma(n) \geq i \\ \sigma[x \rightarrow 2^{\sigma(n)} \cdot \sigma(x), n \rightarrow 0] & 0 < \sigma(n) < i \\ \sigma & \sigma(n) \leq 0 \end{cases}$$


c) An explicit definition for $\bigsqcup_i F^i(\lambda\sigma. \perp)$.

$$\bigsqcup_i F^i(\lambda\sigma. \perp) = \lambda\sigma. \begin{cases} \sigma[x \rightarrow 2^{\sigma(n)} \cdot \sigma(x), n \rightarrow 0] & \sigma(n) > 0 \\ \sigma & \sigma(n) \leq 0 \end{cases}$$

d) Give the least fixed point of F and justify it.

Let $(State \rightarrow State_\perp, \sqsubseteq)$ be a pair of a set and an order relation, we know that $(State \rightarrow State_\perp, \sqsubseteq)$ is a ccpo (proposition pag. 110), and let $F : (State \rightarrow State_\perp) \rightarrow (State \rightarrow State_\perp)$ be a function with $F(f) = cond(\mathcal{B}[b], f \circ S[S], id)$ with $S \in Stm$ and $b \in BExp$, we also know F is monotone and continuous in $(State \rightarrow State_\perp, \sqsubseteq)$ (result in pag 116).

Applying Fixed-Point theorem we conclude that $\bigsqcup_i F^i(\lambda\sigma. \perp)$ is the least fixed point of the function F , i.e

$$lfp F = \bigsqcup_i F^i(\lambda\sigma. \perp) = \lambda\sigma. \begin{cases} \sigma[x \rightarrow 2^{\sigma(n)} \cdot \sigma(x), n \rightarrow 0] & \sigma(n) > 0 \\ \sigma & \sigma(n) \leq 0 \end{cases}$$


d) Execute the program of d) under the given initial state.

Let us consider the following statement

$$S = (x := 1; \text{while } n > 0 \text{ do } (x := 2 * x; n := n - 1))$$

and the initial state $\sigma = [n \rightarrow 4]$.

In order to compute S under σ we have to compose σ with $\mathcal{S}[[S]]$, i.e.

$$\sigma \circ \mathcal{S}[(x := 1; \text{while } n > 0 \text{ do } (x := 2 * x; n := n - 1))]$$

We know the semantics of a sequence $S_1; S_2$ is defined as follows:

$$\mathcal{S}[S_1; S_2] = \mathcal{S}[S_2] \circ \mathcal{S}[S_1]$$

Therefore, we obtain that

$$\mathcal{S}[x := 1; \text{while } n > 0 \text{ do } (x := 2 * x; n := n - 1)] = \mathcal{S}[\text{while } n > 0 \text{ do } (x := 2 * x; n := n - 1)] \circ \mathcal{S}[x := 1]$$

And finally

$$\begin{aligned} \sigma \circ \mathcal{S}[S] &= \sigma \circ \mathcal{S}[(x := 1; \text{while } n > 0 \text{ do } (x := 2 * x; n := n - 1))] \\ &= \sigma \circ \mathcal{S}[\text{while } n > 0 \text{ do } (x := 2 * x; n := n - 1)] \circ \mathcal{S}[x := 1] \\ &= \sigma \circ \mathcal{S}[\text{while } n > 0 \text{ do } (x := 2 * x; n := n - 1)] [x \rightarrow 1] \\ &= [n \rightarrow 4] \circ \mathcal{S}[\text{while } n > 0 \text{ do } (x := 2 * x; n := n - 1)] [x \rightarrow 1] \\ &= \mathcal{S}[\text{while } n > 0 \text{ do } (x := 2 * x; n := n - 1)] [x \rightarrow 1][n \rightarrow 4] \\ &= \mathcal{S}[\text{while } n > 0 \text{ do } (x := 2 * x; n := n - 1)] [x \rightarrow 1, n \rightarrow 4] \\ &= (lfp F) [x \rightarrow 1, n \rightarrow 4] \\ &= \left(\bigsqcup_i F^i(\lambda\sigma. \perp) \right) [x \rightarrow 1, n \rightarrow 4] \\ &= \lambda\sigma'. \sigma' [x \rightarrow 2^4, n \rightarrow 0] \\ &= \lambda\sigma'. \sigma' [x \rightarrow 16, n \rightarrow 0] \end{aligned}$$
