

# Assignments on Program Semantics

Theory of Programming Languages  
Master's Degree in Formal Methods in Computer Science  
Year 2021–2022

## Submission deadlines:

- October 20th: Exercise 1.
- October 27th: Exercises 2 and 3.
- November 4th: Exercise 4.

**Submission instructions:** For each submission, students are required to upload a PDF file with the solution(s) to the exercise(s). You can use  $\text{\LaTeX}$  (recommended) or any other word processor/typesetter. Handwritten submissions are also accepted but, in the latter case, the student has to scan their submission in order to obtain a PDF file. Submissions have to be written in English.

1. Let us assume  $e, e' \in \mathbf{AExp}$  and  $x \in \mathbf{Var}$ . The notation  $e[x/e']$  denotes the result of replacing all occurrences of  $x$  in  $e$  by  $e'$ . For example:  $(x + y)[x/(3 * z)] = (3 * z) + y$ .

(a) Define  $e[x/e']$  in a compositional way.

(b) Prove the following *substitution lemma*: for all  $e, e' \in \mathbf{AExp}$ ,  $x \in \mathbf{Var}$ ,  $\sigma \in \mathbf{State}$ :

$$\mathcal{A}[e[x/e']] \sigma = \mathcal{A}[e] \sigma[x \mapsto \mathcal{A}[e'] \sigma]$$

2. Assume we extend the syntax of *While* statements with a new construct: *repeat S until b*. This statement is executed as follows:

(1) Execute *S*.

(2) Check whether *b* is false. In this case, step back to (1). Otherwise, finish.

Define the big-step and small-step semantic rules for this new construct. You cannot rely on the rules of *while* to define the rules of *repeat*. Finally, prove that *repeat S until b* is equivalent to  $(S; \text{while } \neg b \text{ do } S)$

3. Add the following iterative construct to *While*: *for  $x := e_1$  to  $e_2$  do S*. Define its big-step and small-step semantic rules. You cannot rely on the *while* or *repeat* construct to do this exercise.
4. Given the function  $F : (\mathbf{State} \rightarrow \mathbf{State}_\perp) \rightarrow (\mathbf{State} \rightarrow \mathbf{State}_\perp)$  defined as follows:

$$F(f) = \text{cond}(\mathcal{B}[n > 0], f \circ \mathcal{S}[x := 2 * x; n := n - 1], id)$$

- (a) Give an explicit definition for  $F(\lambda\sigma.\perp)$ ,  $F^2(\lambda\sigma.\perp)$  and  $F^3(\lambda\sigma.\perp)$ .
- (b) From the results above, conjecture a general definition for  $F^i(\lambda\sigma.\perp)$  where  $i \geq 1$ .  
[Optional] Prove by induction on  $i$  that your conjecture is correct.
- (c) Give an explicit definition for  $\bigsqcup_i F^i(\lambda\sigma.\perp)$ .
- (d) Which is the least fixed point of  $F$ ? Justify your answer.
- (e) Given the above, compute the state resulting from the execution of the following program

$x := 1; \text{ while } n > 0 \text{ do } (x := 2 * x; n := n - 1)$

under the initial state  $\sigma = [n \mapsto 4]$ .