

The (simply-typed) λ -calculus

Theory for Programming Languages
MASTER IN FORMAL METHODS FOR SW ENGINEERING
UCM / UPM / UAM

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Exercise 1. Give a proof in TA_λ for the following type assignment for function composition:

$$\vdash \lambda f. \lambda g. \lambda x. f(g x) : (b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c)$$

Exercise 2. Provide a type for the combinator $S = \lambda x. \lambda y. \lambda z. (x z)(y z)$.

Exercise 3. Provide a type for the term $I I = (\lambda x. x)(\lambda x. x)$.

Exercise 4. Provide types for the terms in exercise 6 of the previous sheet (Church booleans) along with their corresponding derivation.

Exercise 5. Prove that if $\Gamma \vdash M : \sigma$ then $\Gamma[\rho \mapsto \tau] \vdash M : \sigma[\rho \mapsto \tau]$.

Exercise 6. Prove that if $\Gamma, x : \sigma \vdash M : \tau$ and $\Gamma \vdash N : \sigma$ then $\Gamma \vdash M[x \mapsto N] : \tau$.

Exercise 7. Prove that if $M \rightsquigarrow_\beta^* M'$ and $\Gamma \vdash M : \sigma$ then $\Gamma \vdash M' : \sigma$.

Ex 1.

$$\begin{array}{c} [x: a] \quad [g: a \longrightarrow b] \\ \hline \frac{}{x: a, \quad g: a \longrightarrow b \quad \vdash g x : b} \text{App} \\ \hline \frac{}{x: a, \quad g: a \longrightarrow b, \quad f: b \longrightarrow c \quad \vdash f(g x) : c} \text{App} \\ \hline \frac{}{x: a, \quad g: a \longrightarrow b, \quad f: b \longrightarrow c \quad \vdash \lambda x. f(g x) : a \longrightarrow c} \text{Abs} \\ \hline \frac{}{x: a, \quad g: a \longrightarrow b, \quad f: b \longrightarrow c \quad \vdash \lambda g. \lambda x. f(g x) : (a \longrightarrow b) \longrightarrow (a \longrightarrow c)} \text{Abs} \\ \hline \frac{}{x: a, \quad g: a \longrightarrow b, \quad f: b \longrightarrow c \quad \vdash \lambda f. \lambda g. \lambda x. f(g x) : (b \longrightarrow c) \longrightarrow (a \longrightarrow b) \longrightarrow (a \longrightarrow c)} \text{Abs} \end{array}$$

Ex 2.

$$S = \lambda x. \lambda y. \lambda z. (xz)(yz)$$

$$\frac{[z:a] \quad [x:a \rightarrow b \rightarrow c]}{z:a, x:a \rightarrow b \rightarrow c \vdash xz:b \rightarrow c} \text{App}$$

$$\frac{[z:a] \quad [y:a \rightarrow b]}{z:a, y:a \rightarrow b \vdash yz:b} \text{App}$$

$$z:a, x:a \rightarrow b \rightarrow c, y:a \rightarrow b \vdash (xz)(yz):c$$

$$\frac{}{z:a, x:a \rightarrow b \rightarrow c, y:a \rightarrow b \vdash \lambda z. (xz)(yz):a \rightarrow c} \text{Abs}$$

$$\frac{}{z:a, x:a \rightarrow b \rightarrow c, y:a \rightarrow b \vdash \lambda y. \lambda z. (xz)(yz):(a \rightarrow b) \rightarrow (a \rightarrow c)} \text{Abs}$$

$$\frac{}{z:a, x:a \rightarrow b \rightarrow c, y:a \rightarrow b \vdash \lambda x. \lambda y. \lambda z. (xz)(yz):(a \rightarrow b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c)} \text{Abs}$$

Ex 3.

$$(\lambda x. x)(\lambda x. x)$$

First of all, we must do α -conversion.

$$(\lambda x. x)(\lambda x. x) \sim_{\alpha} (\lambda x. x)(\lambda y. y)$$

$$\frac{[x:a \rightarrow a] \quad [x:a \rightarrow a]}{x:a \rightarrow a \vdash \lambda x. x:(a \rightarrow a) \rightarrow (a \rightarrow a)} \text{Abs}$$

$$\frac{}{x:a \rightarrow a, y:a \vdash (\lambda x. x)(\lambda y. y):a \rightarrow a} \text{App}$$

$$\frac{[y:a] \quad [y:a]}{y:a \vdash \lambda y. y:a \rightarrow a} \text{Abs}$$

Ex 4.

$$\text{TRUE} = \lambda x. \lambda y. x$$

$$\frac{[x:a] \quad [y:b]}{x:a, y:b \vdash \lambda y. x:b \rightarrow a} \text{Abs}$$

$$\frac{}{x:a, y:b \vdash \lambda x. \lambda y. x:a \rightarrow b \rightarrow a} \text{Abs}$$

$$\text{FALSE} = \lambda x. \lambda y. y$$

$$\frac{[y:b]}{y:b \vdash \lambda y. y:b \rightarrow b} \text{Abs}$$

$$\frac{}{x:a, y:b \vdash \lambda x. \lambda y. y:a \rightarrow b \rightarrow b} \text{Abs}$$

$$\text{COND} = \lambda b. \lambda x. \lambda y. b \times y$$

$$\begin{array}{c} [b: a \rightarrow f \rightarrow c] \quad [x:a] \\ \hline \frac{b: a \rightarrow f \rightarrow c, x:a \vdash bx:f \rightarrow c \quad \text{App}}{b: a \rightarrow f \rightarrow c, x:a, y:f \vdash bx.y:c} \quad [y:b] \\ \hline \frac{b: a \rightarrow f \rightarrow c, x:a, y:f \vdash \lambda y. bxy : f \rightarrow c \quad \text{Abs}}{b: a \rightarrow f \rightarrow c, x:a, y:f \vdash \lambda x. \lambda y. bxy : a \rightarrow f \rightarrow c} \quad \text{Abs} \\ \hline \frac{b: a \rightarrow f \rightarrow c, x:a, y:f \vdash \lambda b. \lambda x. \lambda y. bxy : (a \rightarrow f \rightarrow c) \rightarrow a \rightarrow f \rightarrow c \quad \text{Abs}}{} \end{array}$$

$$\text{CONJ} = \lambda b_1. \lambda b_2. b_1 \cdot b_2 \text{ FALSE}$$

For the proof tree above, we know that $\text{FALSE} : a \rightarrow b \rightarrow b$

$$\begin{array}{c} [b_1: a \rightarrow (a \rightarrow b \rightarrow b) \rightarrow \gamma] \quad [b_2: a] \\ \hline \frac{b_1: a \rightarrow (a \rightarrow b \rightarrow b) \rightarrow \gamma, b_2: a \vdash b_1.b_2: (a \rightarrow b \rightarrow b) \rightarrow \gamma \quad \text{App}}{b_1: a \rightarrow (a \rightarrow b \rightarrow b) \rightarrow \gamma, b_2: a, \text{FALSE} : a \rightarrow b \rightarrow b \vdash b_1.b_2 \text{FALSE} : \gamma} \quad [\text{FALSE} : a \rightarrow b \rightarrow b] \\ \hline \frac{b_1: a \rightarrow (a \rightarrow b \rightarrow b) \rightarrow \gamma, b_2: a, \text{FALSE} : a \rightarrow b \rightarrow b \vdash \lambda b_2. b_1.b_2 \text{FALSE} : a \rightarrow \gamma \quad \text{Abs}}{b_1: a \rightarrow (a \rightarrow b \rightarrow b) \rightarrow \gamma, b_2: a, \text{FALSE} : a \rightarrow b \rightarrow b \vdash \lambda b_1. \lambda b_2. b_1.b_2 \text{FALSE} : (a \rightarrow (a \rightarrow b \rightarrow b) \rightarrow \gamma) \rightarrow a \rightarrow \gamma \quad \text{Abs}} \end{array}$$

$$\text{DISJ} = \lambda b_1. \lambda b_2. b_1 \text{ TRUE} \cdot b_2$$

For the proof tree above, we know that $\text{TRUE} : a \rightarrow b \rightarrow a$

$$\begin{array}{c} [b_1: (a \rightarrow b \rightarrow a) \rightarrow a \rightarrow \gamma] \quad [\text{TRUE} : a \rightarrow b \rightarrow a] \\ \hline \frac{b_1: (a \rightarrow b \rightarrow a) \rightarrow a \rightarrow \gamma, \text{TRUE} : a \rightarrow b \rightarrow a \vdash b_1.\text{TRUE} : a \rightarrow \gamma \quad \text{App}}{b_1: (a \rightarrow b \rightarrow a) \rightarrow a \rightarrow \gamma, \text{TRUE} : a \rightarrow b \rightarrow a, b_2: a \vdash b_1.\text{TRUE}.b_2 : \gamma} \quad [b_2: a] \\ \hline \frac{b_1: (a \rightarrow b \rightarrow a) \rightarrow a \rightarrow \gamma, \text{TRUE} : a \rightarrow b \rightarrow a, b_2: a \vdash \lambda b_2. b_1.\text{TRUE}.b_2 : a \rightarrow \gamma \quad \text{Abs}}{b_1: (a \rightarrow b \rightarrow a) \rightarrow a \rightarrow \gamma, \text{TRUE} : a \rightarrow b \rightarrow a, b_2: a \vdash \lambda b_1. \lambda b_2. b_1.\text{TRUE}.b_2 : ((a \rightarrow b \rightarrow a) \rightarrow a \rightarrow \gamma) \rightarrow a \rightarrow \gamma \quad \text{Abs}} \end{array}$$

$$\text{NEG} = \lambda b. \lambda x. \lambda y. b \ y \ x$$

$$\begin{array}{c} [b: f \rightarrow a \rightarrow c] \quad [y:f] \\ \hline \frac{b: f \rightarrow a \rightarrow c, y:f \vdash by : a \rightarrow c \quad \text{App}}{b: f \rightarrow a \rightarrow c, y:f, x:a \vdash by.x : c} \quad [x:a] \\ \hline \frac{b: f \rightarrow a \rightarrow c, y:f, x:a \vdash \lambda y. byx : f \rightarrow c \quad \text{Abs}}{b: f \rightarrow a \rightarrow c, y:f, x:a \vdash \lambda x. \lambda y. byx : a \rightarrow f \rightarrow c \quad \text{Abs}} \\ \hline \frac{b: f \rightarrow a \rightarrow c, y:f, x:a \vdash \lambda b. \lambda x. \lambda y. byx : (f \rightarrow a \rightarrow c) \rightarrow a \rightarrow f \rightarrow c \quad \text{Abs}}{} \end{array}$$