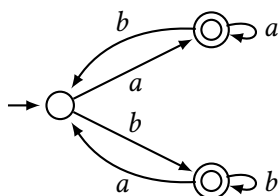


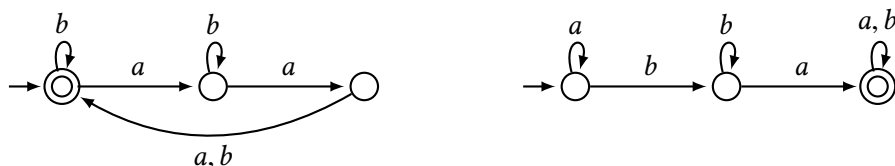
CS 321: Study guide

- True or false? If true, justify. If false, give counterexample:
 - if A not regular then \overline{A} is not regular either.
 - if A is regular and $A \cap B$ is not regular, then B is not regular
 - If A is not regular and B is not regular, then $A \cap B$ is not regular
 - If A is regular and $A \subseteq B$, then B is regular.
 - $a(aa)^* + (aa)^*$ and a^* are equivalent regular expressions
- Show that the following are regular. This can involve describing an explicit DFA, NFA, or regular expression; it can also involve using known closure properties:
 - $\{w \in \{a, b\}^* \mid w \text{ contains } abb \text{ as a substring but not } abba\}$
 - $\{w \in \{a, b\}^* \mid w \text{ contains at least 4 } b\text{'s}\}$
 - $\{w \in \{a, b\}^* \mid w \text{ contains at most 4 } b\text{'s}\}$
 - $\{w \in \{0, 1\}^* \mid w \text{ is a multiple of 5 in binary}\}$
 - $\{w \in \{0, 1\}^* \mid \text{first two characters of } w \text{ equal the last two characters of } w\}$.

For example: abbbbab is in the language, but aaaaabb is not.
- Regarding this NFA . . .



- Convert it to an equivalent DFA
 - Convert it to an equivalent regular expression
- Draw a DFA that accepts the intersection of these two languages:



- Convert $(aa^*b + ba^*b)^*ba^*$ to an NFA
- Show that the following languages are **not regular**. This can involve the pumping lemma and/or closure properties.
The following box **will be included** on the exam. You don't have to memorize the pumping-lemma game:

For reference, here is the *pumping lemma game* (for language A):

1. Adversary picks a number $p \geq 0$.
2. You pick a string $w \in A$, such that $|w| \geq p$.
3. Adversary breaks w into $w = xyz$, such that $|xy| \leq p$ and $y \neq \epsilon$.
4. You pick a number $i \geq 0$. If $xy^iz \notin A$, then you win.

If you can describe a strategy in which you always win, then A is not regular.

- (a) $\{a^m b^n a^n b^m \mid m \geq 0 \text{ and } n \geq 0\}$
- (b) $\{a^m b^n c^k \mid m < n < k\}$
- (c) $\{a^m b^n c^k \mid m, n, k \text{ are all distinct}\}$
- (d) $\{a^m b^n \mid m < n \text{ and both } n \text{ and } m \text{ are multiples of } 3\}$
- (e) $\{w \in \{[, \cdot, \cdot, \cdot]\}^* \mid w \text{ is a string of balanced parentheses}\}$