

Data Engineering

1) SQL

→ Case Example. (Can count ^{multiple} values for diff columns using single query.)

```
select sum(case col1
              when 2.99 then 1
              else 0
              END) as regular,
       sum(case col2
              when 4.99 then 1
              else 0
              END) as rental rate;
```

→ COALESCE(Function).

→ Accepts unlimited no. of arguments and returns 1st argument i.e not null if all null returns null.

Eg! - SELECT coalesce(NULL, 1, 2, 3) → 1

★ → We use COALESCE to substitute null values while doing operations.

Eg! - SELECT Item, (price - COALESCE(discount, 0))

As final-price from table.

Replace discount with value null to 0

~~Cast~~ Cast (operator) :- It lets you convert from one data type to other.

Syntax :- `SELECT CAST('5' AS INTEGER)`
Postgres SQL \Rightarrow `SELECT '5'::INTEGER.`

\rightarrow NULLIF (function) :- Returns null if 2 args passed are '=' else returns 1st arg passed.

\rightarrow View :- A database obj that is of stored query.
 \rightarrow Adv from table is no phy storage and data queried is up to date.

Eg :- `CREATE VIEW view_name AS`
`SELECT c1, c2 FROM t1 inner join address`
`ON c1.a-id = c2.a-id.`

To update view :-

`CREATE OR REPLACE view_name AS`
`query.....`

Drop view if exists :-

`DROP VIEW IF EXISTS view_name.`

Rename view :-

`ALTER VIEW v1 RENAME to to v2`

\rightarrow Importing and exporting data functionality of pg Admin
that allows us to import data from .CSV file to already existing table.

→ psycopg2:- Python library to interact with DB in PostgreSQL.

Eg:- import psycopg2 as pg2
conn = pg2.connect(database='name', user='postgres', password='pwd')
cur = conn.cursor()
cur.execute('SELECT * FROM t1')
cur.fetchone() ⇒ Fetches query results
cur.fetchmany(10) ⇒
cur.fetchall() ⇒
Query = """CREATE TABLE - - - - -"""
cur.execute(query).
★ cur.commit() ⇒ To commit changes to DB.
★ conn.close()

Window Functions:- They perform aggregate ops on grp of rows but they produce result for each row. Eg:- SELECT c1, c2, c3, avg(c3) over(partition by c2)

as dept c3_avg from t1. (group by c2)

Eg:- SELECT c1, avg(c2) over() from emp;

Eg:- select emp.no, dept, sal, count(*) over (PARTITION BY dept) as dept_count from emp;

*Eg 4:- SELECT emp-no, dept, sal,
 SUM(salary) over (PARTITION BY department
 ORDER BY SALARY) <sup>or rolling-dept-
 salary</sup>
 MIN(salary) over (" DESC ")

*Eg 5:- SELECT emp-no, dept, salary,
 RANK() over (order by ~~dept~~ ^{salary} DESC).

RANK() ⇒ gives rank based on column value.

*L> It might miss a rank if there are
 3 people with rank 1 it will miss rank 2, 3 the next
 rank will be 4.

DENSE-RANK() → Works like RANK() but
 ranks will not miss a number if even ^{preceding} ranks
 are same.

Eg 6:- SELECT emp-no, dept, salary
 NTILE(6) over (order by salary DESC).

Eg 7:- (FIRST-VALUE), (LAST-VALUE), (NTH-VALUE).

Eg 8:- (LAG, LEAD), LAG(col1, 2)

★ → We can create table ~~us~~ from query.

Materialized View:-

→ To avoid having redundant storage and up to date data we use views but for complex queries to execute on view takes time so we use materialized views.

★ → Data is stored physically. ~~Q~~

⇒ Gg! CREATE MATERIALIZED VIEW v-name
AS query.

Gg! REFRESH MATERIALIZED VIEW

(To update view with new data from the view query)

→ We can update data manually on the triggers.

★ Gg! ALTER MATERIALIZED VIEW
DROP MATERIALIZED VIEW

(named subquery)
Common table expr (CTE):-

→ To use query result in other query.

★ Syntax:- WITH cte-name AS

SELECT - - - - -)

SELECT - - - - - cte-name - - -

Multi CTE

Syntax with name as (),
name as ()
Select - - -

→ We can not use aliases for columns in same query level.

Indexing:-

→ Single level (primary, clustered, secondary).

→ Multi level (BTree, B+tree)

* It reduces time take to (data transfer b/w RAM and hard disk).

→ It reduces I/O cost.

→ We use sparse when data is sorted

Indexing types.

→ Here ~~only~~ the whole block pointer is stored.

→ We use Dense when data/records is unsorted.

→ Here the pointer is each row pointer is stored.

→ Primary Indexing:-

→ Ordered file where records (key+pointer) of fixed length.

→ Clustered Indexing:-

→ Records are physically ordered on non-key field which doesn't have distinct value for each record

~~for each block of disti~~

→ Each distinct value of non-key attribute an entry is created.

→ ~~Each~~ ^{for} Duplicate non-key attribute ~~points to~~ single pointer is used for reference.

Secondary Indexing:-

→ Creating index on ~~sec~~^{2nd} attribute when we have primary indexing.

→ 2nd attr is not ordered in DB. But can order it in index table.

1^o index syntax:-

→ CREATE INDEX index-name

on table-name (col-name);

→ DROP INDEX index-name;
(col 1, col 2); multi ~~col~~ index.

→ Indexing should be reconsidered when:-

→ Small tables

→ Have frequent large batch update, insert ops.

→ On columns with high no. of null values.

→ Columns that are frequently manipulated.

→ Execute sp-helppindex table-name; ^{* to see all the indexes} ⇒ on particular table.

→ In SQL server
 ↗ Clustered index.
 ↘ Non-clustered index.

→ Clustered Index:-

→ Defines the order in which data is physically stored in a table. As ~~extra~~ sorting can be done in only way we have one clustered index per table. In SQL server 1^o key constraint auto creates

clustered index.

→ When new data is added it maintains order based on cluster index value.

→ Custom clustered index:-

CREATE CLUSTERED INDEX index-name
on table-name (col1, col2);

→ Non clustered index:-

→ Separate storage is needed.

→ Can have multiple non-clustered indexes per table.

Types of DB:- 1. Key-value, 2. Wide column, (used when freq writes and not much read and update).
(Redis, memcache) (Cassandra, Hbase) (Scalable, CQL)

3. Document DB (MongoDB, Firestore, DynamoDB, CouchDB)

→ Unstructured, No schema.

→ Collection (contain zero or more documents).

→ Document (contain key-value pairs JSON).

* → Field ^{within} a collection can be indexed and collections can be organized into a logical hierarchical

→ Reads are much faster while write and update is complex.

4. RDB:- (MySQL, PostgreSQL, SQL server...).

→ ACID compliant.

→ Cockroach DB scales horizontally without reconfig..

5. Graph DB :- (Neo4J, Dgraph).

- Nodes contain data.
- Edges define relationship

→ Good choice alt for RDB if we have lot of joins and perf issue.

6. Search Engine :- (Give list of most appropriate results).

- Most are on top of Apache Lucene

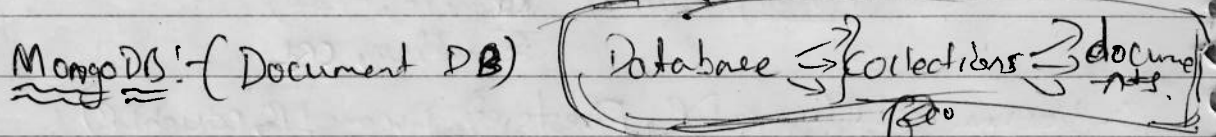
Eg:- Solr, elastic search.

dhbs

7. MultiModel (FAUNA) :-

- Use GraphQL to describe how we want data

MongoDB :- (Document DB)



- Documents can be stored in formats like JSON, BSON, XML.

→ BSON :- Binary JSON to overcome JSON drawback

↳ It encodes type and length information makes it possible to traverse more quick compared to JSON.

↳ It adds some non-JSON data types (date, binary data)

↳ MongoDB driver converts JSON to BSON.

- Some Document DB provide schema validation.

→ Collections :- Set of documents that can have different schema. Eg:- Users collection can have 2 documents with 1 user detail in each doc.

→ CRUD can be done here.

↳ using unique id/field values. Indexes can be used for good perf.

→ Document DB's are distributed. (Horizontal scaling).

→ Single transaction works fine in MongoDB.

→ ~~For mult~~ MongoDB support multi transactions but the perf might be affected.

→ Cluster is place where MongoDB DB's are stored.

★ → Database design and data modeling are major factors in MongoDB performance.

★ → ~~where~~ -id will be created automatically for each new doc created in collection. But we can assign it manually using '-id' as key.

★ → \$ with word in MongoDB indicate its reserved word.

→ db.flightData.updateMany(
 ↓
collection {_id: {}, \$set: {mark: "delete"}},
 ↓
filter set to empty to update all docs

→ db.flightData.insertMany([{}])
 ↑
pass collection documents here.

→ db.flightData.find({distance: {\$gt: 10000}}).pretty().
 ↑
(to get data distance > 10000)

→ db.flightData.update({_id: ObjectId("xxx")}, {del: false})
 This will not just set the del field to false
 it will replace whole obj with only _id and del keys

* \rightarrow `db.pamengers.find()` \Rightarrow Returns cursor to go through elements

\rightarrow `db.pamengers.find().toArray()` \Rightarrow Returns all elements (documents)

\rightarrow `db.pamengers.find().forEach()` \Rightarrow Returns 1 element for a loop reducing load...

* Projection:- Instead of fetching whole document (or) object and filtering it in app which increase load and latency we can filter (or) get the data only we want ~~inside~~ ~~from~~ by filtering within DB we use projection.

Eg:- `db.pamengers.find({}, {name: 1})` \Rightarrow Returns all data with id, name

* `db.pamengers.find({}, {name: 1, -id: 1})` \Rightarrow Returns all data with only name

* Embedded Document:- Document inside other document.

\rightarrow Max document size in MongoDB is 16MB.

\rightarrow `db.hobbies.insertMany([{}], {ordered: false});`

\rightarrow By default ordered is true so if any doc insertion fails

only preceding docs would get inserted to avoid it we can set ~~ordered~~ ~~false~~

\rightarrow We use journal ~~in~~ ~~case~~ to maintain to-do tasks so that in case the memory ~~is~~ get wiped out (power) it still knows what to do during journal. $\{j: \text{undefined (no response if it's not in journal)}, j: \text{True (response)}$

\rightarrow Comment in shell:-

`mongoimport -u shows.json -d movieData -c movies`
need to be specified if not in file. \rightarrow db name \rightarrow collection name
 \rightarrow jsonArray \rightarrow drop
with \rightarrow exist \rightarrow without \rightarrow delete & append create