Homework 10

Solutions

1.
$$Q1* = Q1' + Q2$$

$$Q2* = X.Q2'$$

$$Z = Q1 + Q2$$

| Q1 | Q2 | Q1*Q2* | | 7 | L |
|----|----|--------|------|-----|-----|
| | | X =0 | X =1 | X=0 | X=1 |
| 0 | 0 | 10 | 11 | 1 | 1 |
| 0 | 1 | 10 | 10 | 0 | 0 |
| 1 | 0 | 00 | 01 | 1 | 1 |
| 1 | 1 | 10 | 10 | 1 | 1 |

| Present state | Input X = 0 | Input X = 1 |
|---------------|-------------|-------------|
| | Next sta | te, Output |
| A | C,1 | D,1 |
| В | C,0 | C,0 |
| С | A,1 | B,1 |
| D | C,1 | C,1 |

2. Overlapping sequence 1001

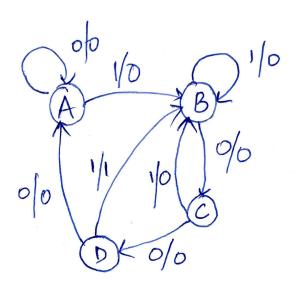
State A - 3 or more 1s

State B - 1 or more consecutive 1s

State C - 10

State D - 100

| Present state | Input $X = 0$ | Input X = 1 |
|---------------|---------------|-------------|
| | Next sta | te, Output |
| A | A, 0 | B,0 |
| В | C,0 | B,0 |
| С | D,0 | B,0 |
| D | A,0 | B,1 |

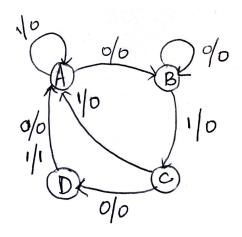


3. State A - 0101/0100/2 or more consequetive 1s

State B - 1 or more 0s after state A

State C - 01

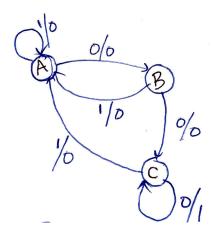
State D - 010



4. State A - 1or more 1s

State B - 0

State C - 2 or more 0s



5.

| Present state | Input X = 0 | Input X = 1 |
|---------------|-------------|-------------|
| | Next sta | te, Output |
| A | B, 0 | C,1 |
| В | D,0 | F,1 |

| С | F,0 | E,0 |
|---|-----|-----|
| D | B,0 | G,1 |
| Е | F,0 | C,0 |
| F | E,0 | D,0 |
| G | F,0 | G,0 |

$$P0 = (ABCDEFG)$$

To perform state minimization, we first separate the states based on the outputs for a 1 bit input.

For X=0, all outputs are same.

For X=1, outputs are 1 for ABD and 0 for CEFG.

Hence,
$$P1 = (ABD)(CEFG)$$

P2:

$$X=0: ABD \Rightarrow BDB$$

CEFG => FFEF

$$X=1: ABD \Rightarrow CFG$$

Since ECG and D are in separate set of states, we can separate F from CEG

Thus
$$P2 = (ABD)(CEG)(F)$$

P3:

$$X=0:ABD \Rightarrow BDB$$

$$CEG \Rightarrow FFF$$

$$F \Rightarrow E$$

$$X=1:ABD \Rightarrow CFG$$

$$CEG \Rightarrow ECG$$

$$F \Rightarrow D$$

Since CG and F are in separate set of states, we can separate B from AD.

Thus
$$P3 = (AD)(B)(CEG)(F)$$

P4:

$$X=0:AD \Rightarrow BB$$

 $B \Rightarrow D$

 $CEG \Rightarrow FFF$

 $F \Rightarrow E$

$$X=1 : AD => CG$$

 $B \Rightarrow F$

 $CEG \Rightarrow ECG$

 $F \Rightarrow D$

$$P4 = (AD)(B)(CEG)(F)$$

Since P3=P4, no further separation can be done,

From P4, it is evident that A, D are equivalent states. And C,E,G are equivalent states.

We can therefore remove D,E,G from the state table.

Thus after minimization the state table looks like:

| Present state | Input X = 0 | Input X = 1 |
|---------------|-------------|-------------|
| | Next sta | te, Output |
| A | B, 0 | C,1 |
| В | A,0 | F,1 |
| С | F,0 | C,0 |
| F | C,0 | A,0 |

| Present state | Input X = 0 | Input X = 1 |
|---------------|-------------|-------------|
| | Next sta | te, Output |
| A | A, 0 | B,0 |
| В | C,0 | B,0 |
| С | D,0 | B,0 |
| D | A,0 | E,0 |
| Е | C,0 | B,1 |

State minimization:

P0 = (ABCDE)

P1:

X=0, all outputs are 0.

X=1, outputs are 0 for ABCD and 1 for E

P1 = (ABCD)(E)

P2:

 $X=0: ABCD \Rightarrow ACDA$

 $E \Rightarrow C$

 $X=1: ABCD \Rightarrow BBBE$

 $E \Rightarrow B$

P2 = (ABC)(D)(E)

P3:

 $X=0: ABC \Rightarrow ACD$

 $D \Rightarrow A$

 $E \Rightarrow C$

 $X=1: ABC \Rightarrow BBB$

 $D \Rightarrow E$

$$E \Rightarrow B$$

$$P3 = (AB)(C)(D)(E)$$

P4:

$$X=0: AB \Rightarrow AC$$

 $C \Rightarrow D$

 $D \Rightarrow A$

 $E \Rightarrow C$

$$X=1: AB => BB$$

 $C \Rightarrow B$

 $D \Rightarrow E$

 $E \Rightarrow B$

$$P4 = (A)(B)(C)(D)(E)$$

No minimization is possible.

State assignment:

| Present state | Input X = 0 | Input X = 1 | |
|---------------|-----------------------------------|-------------|--|
| Q0Q1Q2 | Next state, Output Q0*Q1*Q2*,Z | | |
| 000 | 000, 0 | 001,0 | |
| 001 | 010,0 | 001,0 | |
| 010 | 011,0 | 001,0 | |
| 011 | 000,0 | 100,0 | |
| 100 | 010,0 | 001,1 | |

| Q2X | 00 | 01 | 11 | 10 |
|---------|----|----|----|----|
| Q0Q1 00 | | | | |
| 01 | | | 1 | |
| 11 | X | X | X | X |

| - 1 | | | | |
|-----|----|--|-----|-----|
| - 1 | 10 | | l X | l X |
| - 1 | 10 | | 71 | 71 |

Q0* = Q1Q2X

Q2X

| Q0Q1 00 | | | | 1 |
|---------|---|---|---|---|
| 01 | 1 | | | |
| 11 | X | X | X | X |
| 10 | 1 | | X | X |

Q1* = Q1Q2'X' + Q1'Q2X' + Q0Q2'X'

Q2X

| Q0Q1 00 | | 1 | 1 | |
|---------|---|---|---|---|
| 01 | 1 | 1 | | |
| 11 | X | X | X | X |
| 10 | | 1 | X | X |

Q2* = Q2'X + Q1Q2' + Q0'Q1'X

Q2X

| Q0Q1 00 | | | | |
|---------|---|---|---|---|
| 01 | | | | |
| 11 | X | X | X | X |
| 10 | | 1 | X | X |

Z = Q0X

7. .

| Q0Q1Q2 | Q0*Q1*Q2* D0D1D2 | | |
|--------|---------------------|------|--|
| | UP=1 | UP=0 | |
| 000 | 001 | 111 | |
| 001 | 010 | 000 | |
| 010 | 011 | 001 | |
| 011 | 100 | 010 | |
| 100 | 101 | 011 | |
| 101 | 110 | 100 | |
| 110 | 111 | 101 | |
| 111 | 000 | 110 | |

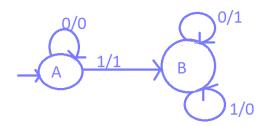
8. Total no. of unique states required = (70+5+75)/5 = 30

No. of flip flops required = n

$$2^n >= 30$$

Thus,
$$n = 5$$

9.



We know the regular method to find 2's complement. Another technique that is useful to draw a state machine is described below:

Start from the right. Do not change the no.s until you reach a 1. Leave the first 1 as such. Complement every no. to the left of the 1.

10. In Moore model the output depends only on the present state.

In Mealy model the output depends on both the present state and the inputs.

```
module FSM(CLOCK,XZ);
input CLOCK,X;
output reg Z;
reg [2:0] Sreg, Snext;
parameter [2:0] S0=3'b000,
S1=3'b001,
S2=3'b010,
S3=3'b011
/* create state memory
always @(posedge CLOCK)
Sreg <= Snext;</pre>
/* create next-state logic */
always @(X,Sreg)
begin
case(Sreg)
S0: if(X==0) Snext=S0;
else Snext=S1;
S1: if(X==0) Snext=S2;
else Snext=S0;
S2: if(X==0) Snext=S3;
else Snext=S2;
S3: if(X==0) Snext=S1;
Snext=S0;
default: Snext=S0;
endcase
end
/* create output logic */
always @(Sreg)
case(Sreg
S0,S1,S2: Z=0;
S3: Z=1;
default: Z=0;
endcase
endmodule
```