

PROBLEM-I

①

Given classes: Side, A, B, C, D

As we have 4 classes A, B, C, D. So the no of bindings for polymorphic testing calls are 4 Bindings.

Hence.

Object Pa Bindings

- * Object of A
- * Object of B
- * Object of C
- * Object of D

Pb Bindings

- * Object of A
- * Object of B
- * Object of C
- * Object of D

Pc Bindings

- * Object of A
- * Object of B
- * Object of C
- * Object of D

Pd Bindings

- * Object of A
- * Object of B
- * Object of C
- * Object of D

Polymorphic Testing at Statement 10: $\text{if}(\text{pa} \rightarrow \text{get}()) > \text{pb} \rightarrow \text{get}())\}$

Pa:

Object of A: Test#1: 12, 13, 14, 0

Object of B: Not Possible

Object of C: Not Possible

Object of D: Not Possible

Pb:

Object of A: ^{Not} Possible

Object of B: Test#2: 12, 13, 14, 0

Object of C: Not Possible

Object of D: Not Possible.

Polymorphic Testing at Statement 14: $\text{if}(\text{pa} \rightarrow \text{get}()) > \text{pc} \rightarrow \text{get}())$

Pa:

Object of A: Test#3: 12, 13, 14, 0

Object of B: Test#4: 13, 12, 14, 0

Object of C: Not Possible

Object of D: Not Possible

Pc:

Object of A: Not Possible

Object of B: Not Possible

Object of C: Test#5: 12, 13, 14, 0

Object of D: Test#6: 12, 13, 14, 0

Polymorphic Testing at Statement 18: $\text{if}(\underline{pb} \rightarrow \text{get}()) > \underline{pc} \rightarrow \text{get}()) \{\}$

pb

Object of A: Test #7: 13, 12, 14, 0

Object of B: Test #8: 12, 13, 14, 0

Object of C: Not Possible

Object of D: Not Possible

pc

Object of A: ~~Test #9: 12, 13, 11, 1~~

Object of B: ~~Test #10: 13, 12, 11, 1~~

Object of C: ~~Test #11: 12, 13, 14, 1~~

Object of D: Test #11: 12, 13, 14, 0.

Polymorphic Testing at Statement 22: $\text{if}(\underline{pa} \rightarrow \text{get}()) + \underline{pb} \rightarrow \text{get}() \leq \underline{pc} \rightarrow \text{get}()$

pa

Object of A: Test #12: 12, 13, 14, 0

Object of B: Test #13: 13, 12, 14, 0

Object of C: Test #14: 12, 13, 11, 1

Object of D: Test #15: 12, 13, 11, 0

pb

Object of A: Test #16: 13, 12, 14, 0

Object of B: Test #17: 12, 13, 14, 0

Object of C: Test #18: 12, 14, 13, 1

Object of D: Test #19: 12, 14, 13, 0

pc

Object of A: Test #20: 12, 12, 11, 1

Object of B: Test #21: 12, 14, 13, 1

Object of C: Test #22: 12, 13, 14, 1

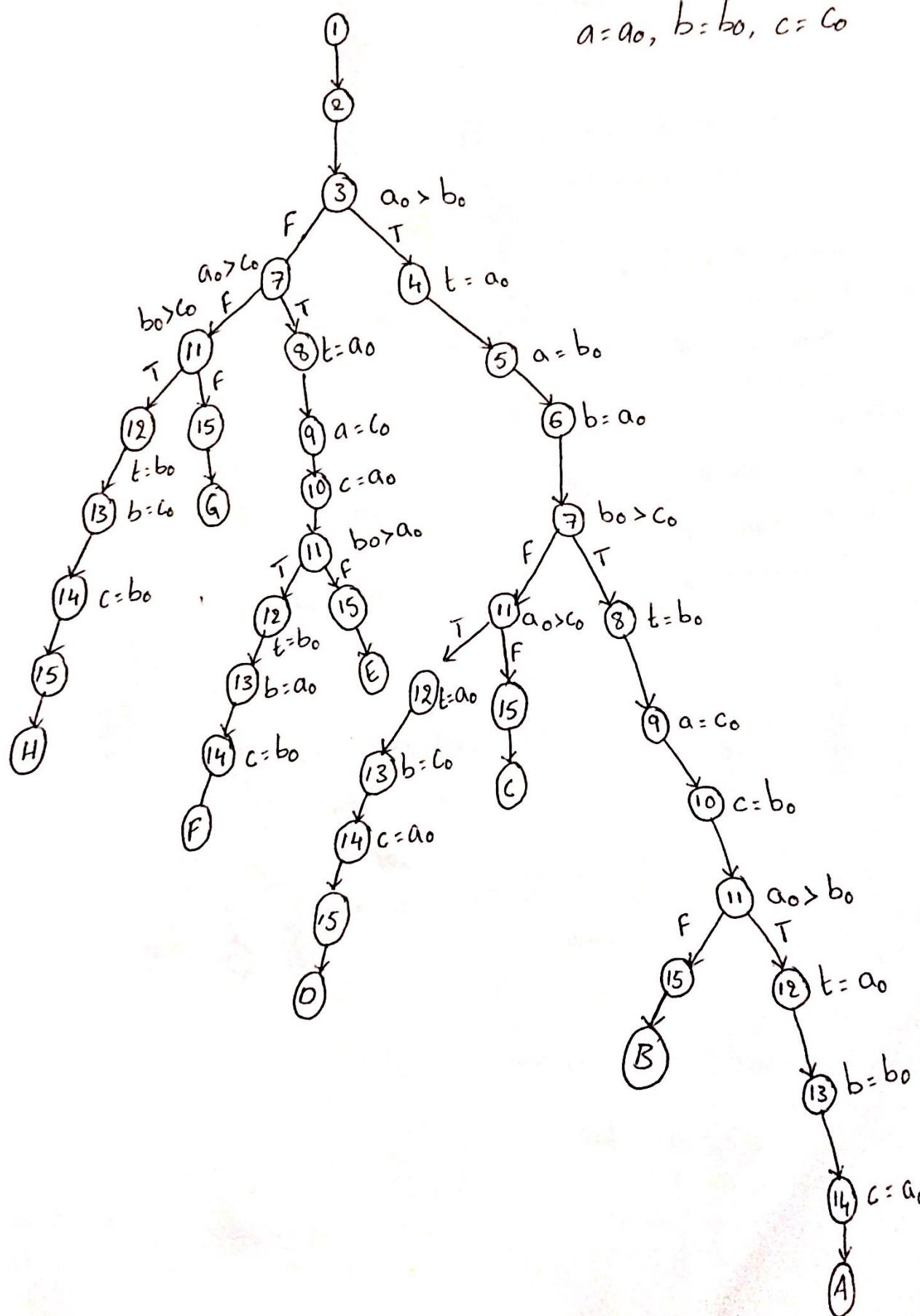
Object of D: Test #23: 12, 13, 14, 0.

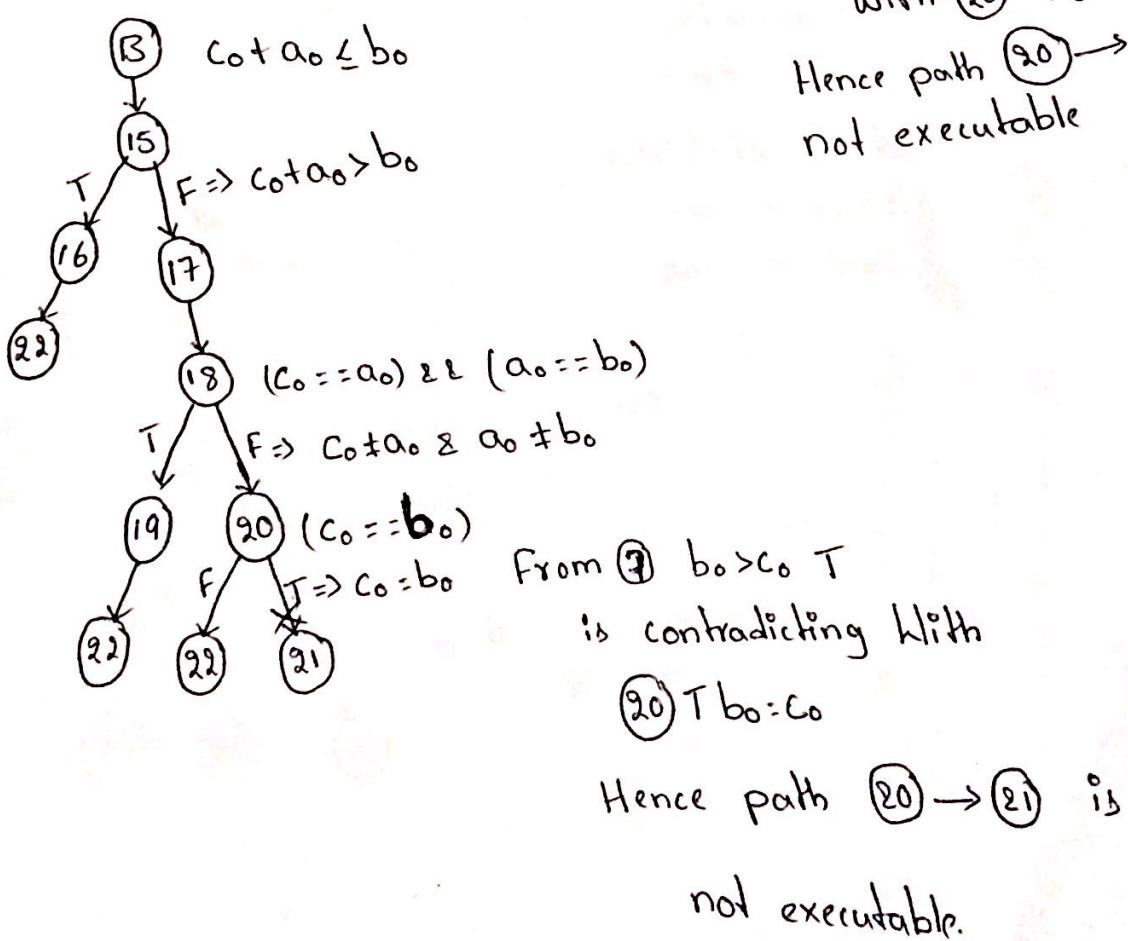
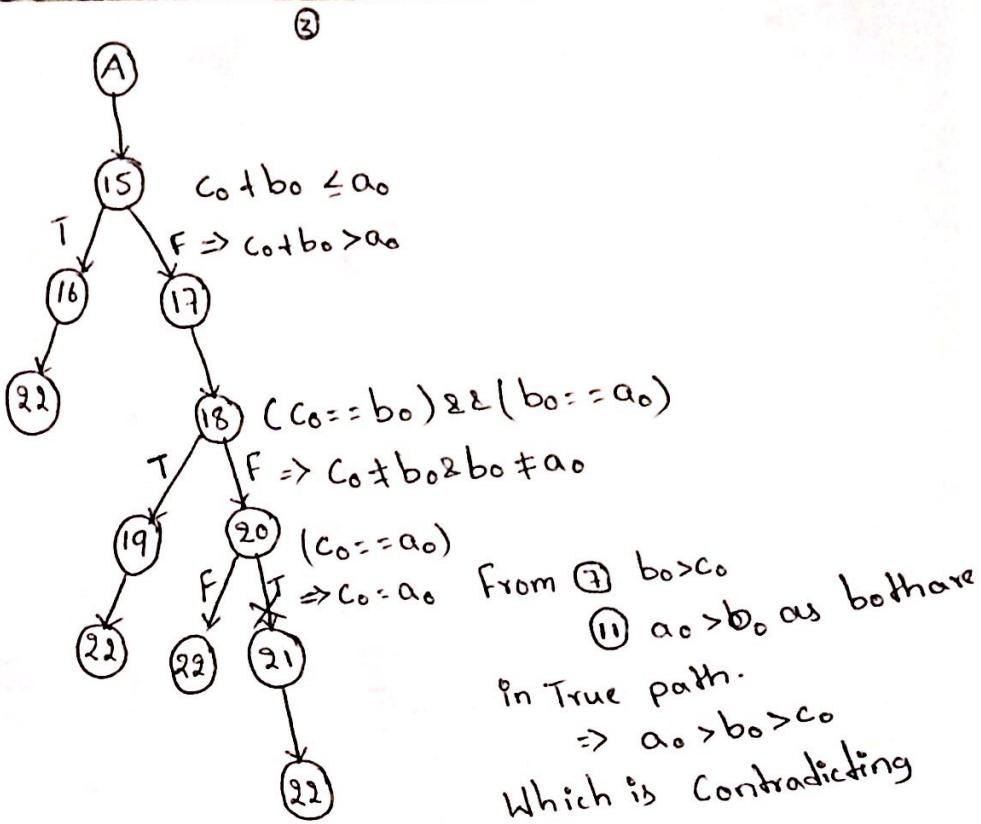
Problem - II

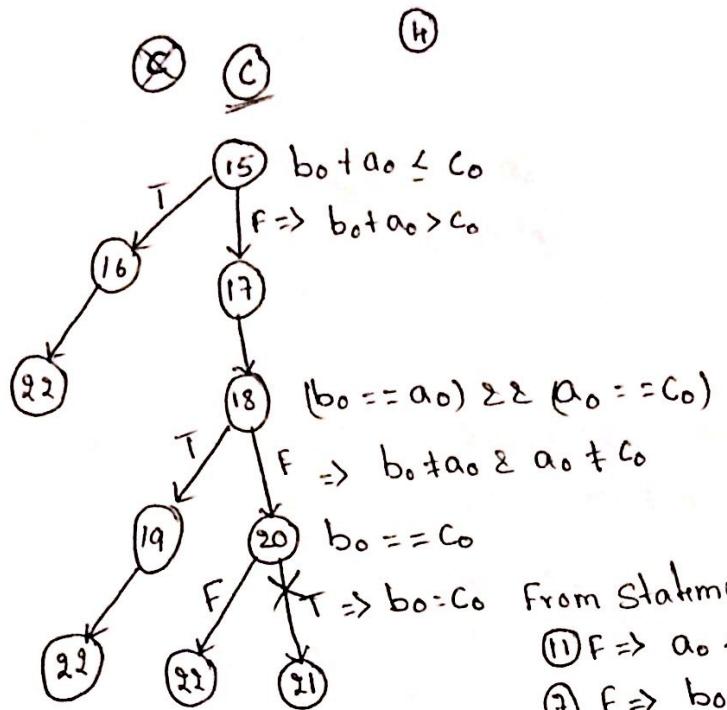
② Symbolic-Execution Tree

Consider

$$a=a_0, b=b_0, c=c_0$$





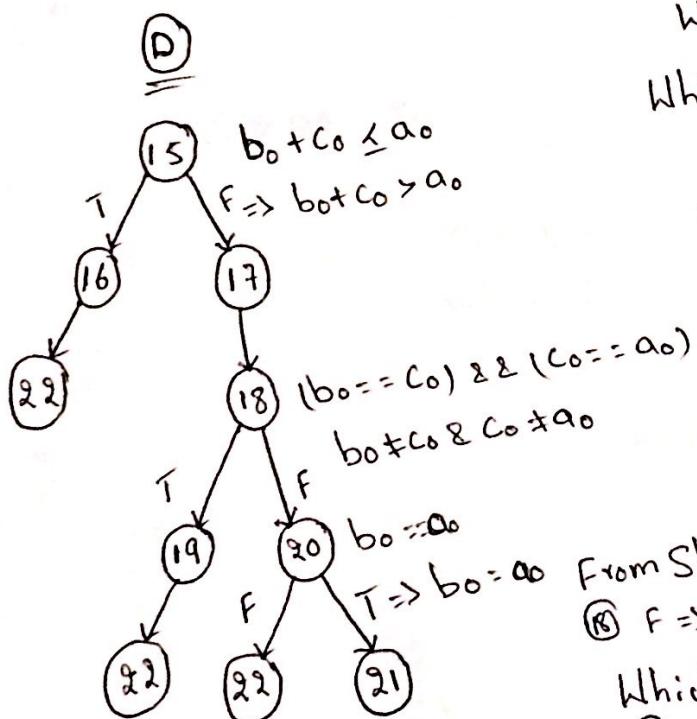


- ⑪ $F \Rightarrow a_o \leq c_o$
- ⑫ $F \Rightarrow b_o \leq c_o$
- ⑬ $T \Rightarrow a_o > b_o$
- ⑭ $F \Rightarrow a_o \neq c_o$

Consolidating all these conditions

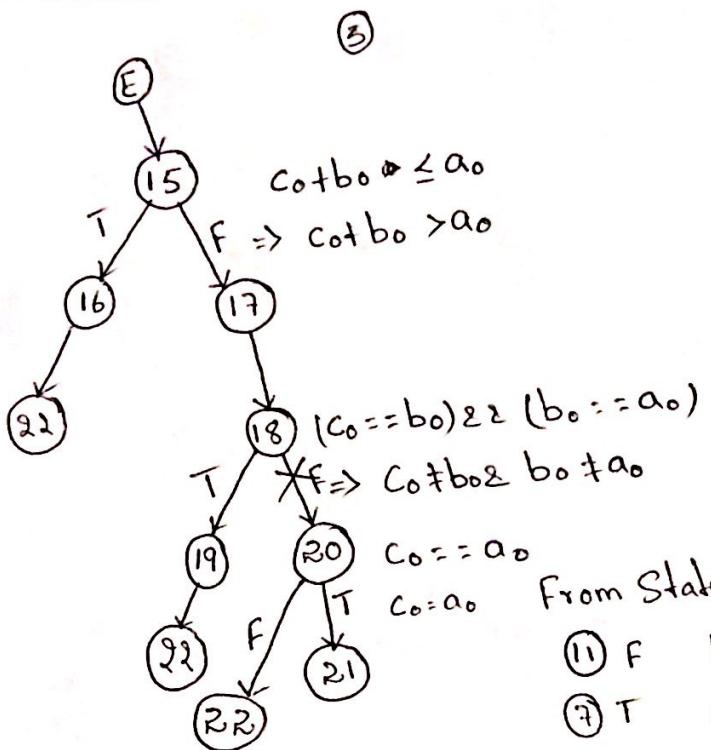
We get. $b_o < a_o \leq c_o$

Which is contradicting with ⑩ T
 $b_o = c_o$. Hence path ⑩ → ⑪ is not executable.



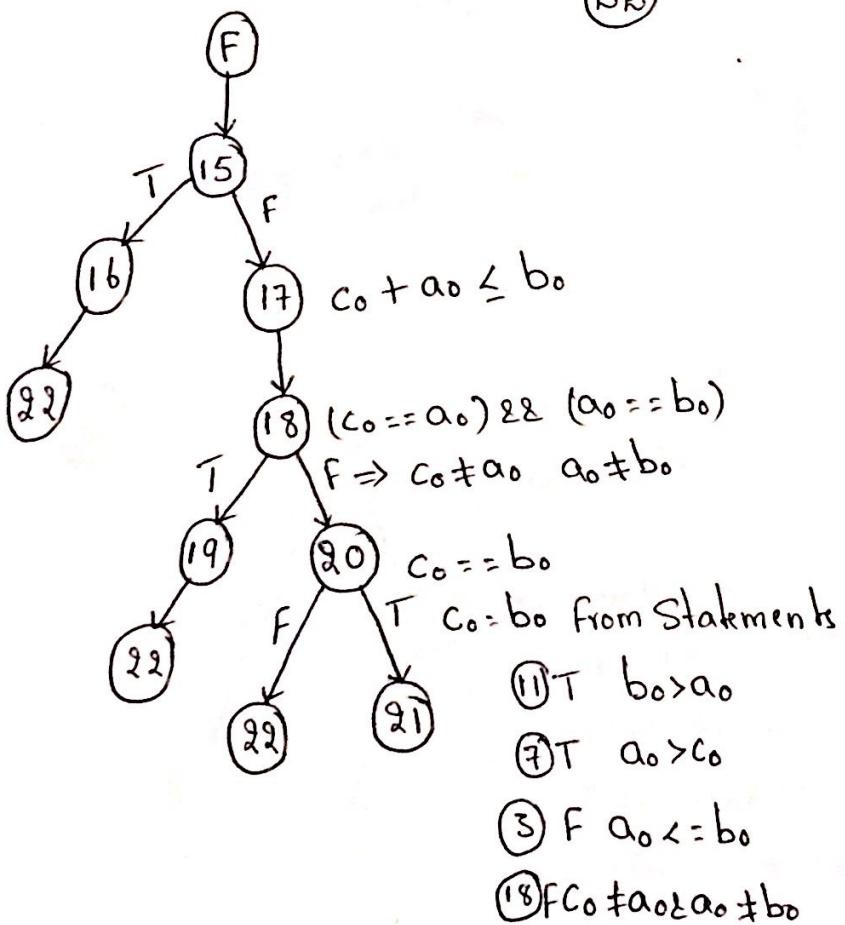
- ⑮ $F \Rightarrow c_o + a_o \geq b_o \neq c_o, ⑯ T \Rightarrow a_o > c_o$
- ⑰ $F \Rightarrow b_o \leq c_o$
- Which is contradicting with ③ $a_o > b_o$ (T)
- ⑩ T for consolidated value

$b_o < a_o \leq c_o$
with $b_o = a_o$,
Hence path ⑩ → ⑪ is not executable.



- From Statement 8:
- (11) F $b_o \leq a_o$
 - (12) T $a_o > c_o$
 - (13) F $a_o \leq b_o$
 - (18) F $c_o \neq a_o \& b_o \neq a_o$

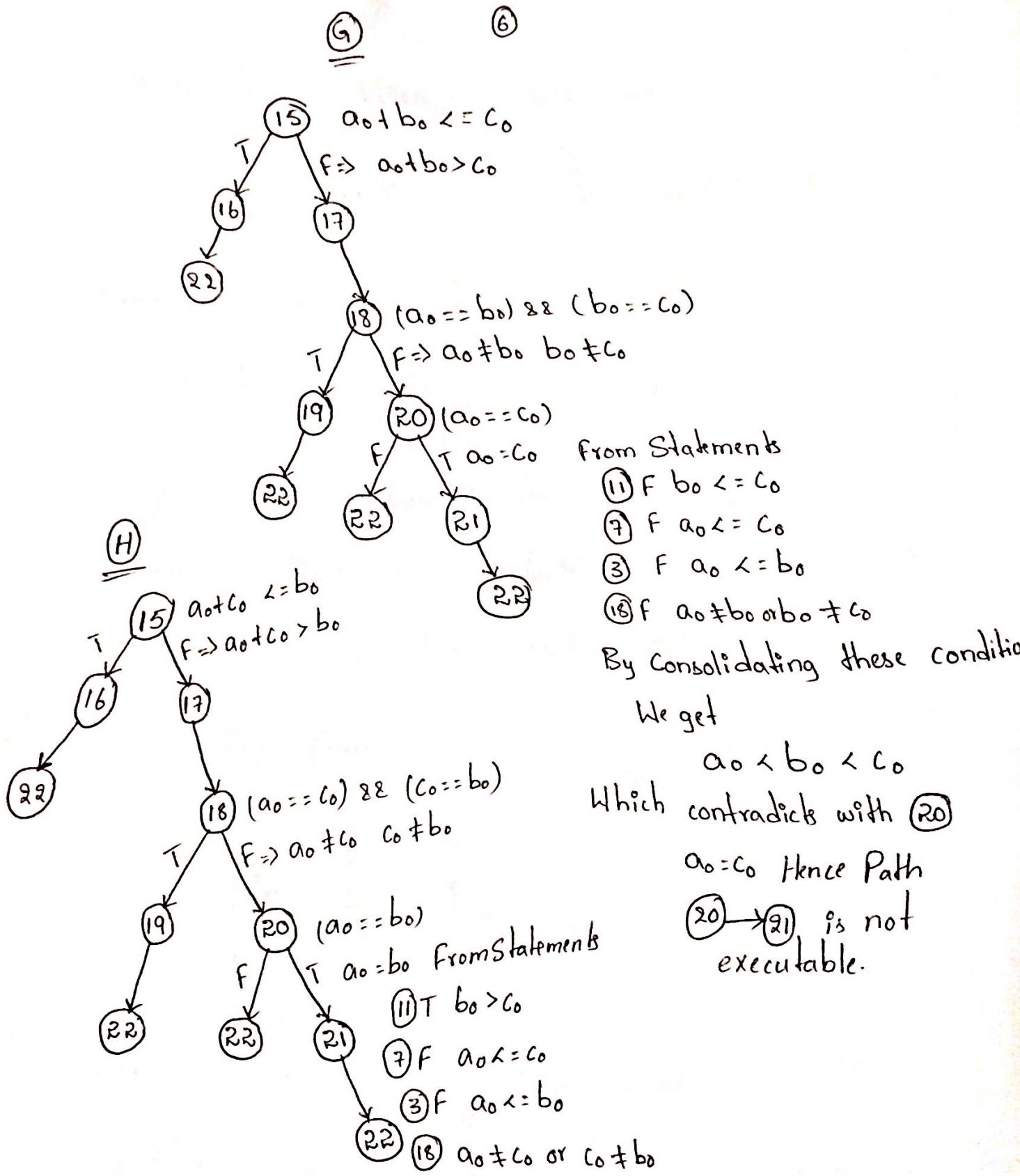
By consolidating all these
We get Contradiction
for (11), (3) and (18)
Statements
i.e; from (11) & (3) we can
infer $a_o = b_o$
But (18) contradicts saying
 $a_o \neq b_o$
Hence this path is not executed



By consolidating We get
 $c_o \leq a_o \leq b_o$

Which contradicts Statement (20) T $c_o = b_o$

Hence this path is not executed.



By consolidating these conditions

We get

$$a_0 < b_0 < c_0$$

Which contradicts with ⑯

$a_0 = c_0$ Hence Path

$\textcircled{20} \rightarrow \textcircled{21}$ is not executable.

By consolidating these Conditions
We get.

$a_0 < c_0 & b_0$ which contradicts with
path $\textcircled{20} \rightarrow \textcircled{21}$ is not $a_0 = b_0$, Hence

Problem - III

(7)

Given Pre-Condition $1 \leq n \leq 100$

Post-Condition $S = \sum_{j=1}^n (|a[j]| == x)$

Loop Invariant at statement 4 is

Consider $S = \sum_{j=i+1}^n (|a[j]| == x)$

Prove the correctness of loop invariant !!!

Using mathematical induction

k : # of times execution reaches point 4.

① $k: 1$ Loop Entry

Path: 1, 2, 3, 4
 $i \downarrow$ \downarrow \downarrow
 $i=n$ $S=0$ $i > 0$

$$\Rightarrow i = n$$

$$S = 0$$

Loop Invariant: $S = \sum_{j=n+1}^n (|a[j]| == x)$

Substituting $i=n$ $\Rightarrow S = \sum_{j=n+1}^n (|a[j]| == x) \rightarrow S=0$ at Step 3

No summation performed hence $S=0$ is true,,

⑧

II Assume that $S_k = \sum_{j=0}^n (|a[j]| = x)$ is true for some k^{th} iteration.

On loop iteration:

We have to show the loop invariant holds true for $(k+1)^{th}$ iteration.

Variables that are modified during iteration is $a[i]$'s

$$\Rightarrow S_{k+1} = \sum_{j=0}^n (|a[j]| = x)$$

To prove this we have to consider the four possible paths available during the iteration.

Path #1: 4, 5, 6, 8, 9, 10, 4

Path #2: 4, 5, 7, 8, 9, 10, 4

Path #3: 4, 5, 6, 8, 10, 4

Path #4: 4, 5, 7, 8, 10, 4.

Case 1: $a[i_k] < 0$ ($-a[i_k] = x$) $\overset{①}{\circ} i_{k+1} = \overset{②}{i}_{k-1}$

Path: $4, \overset{\uparrow}{5}, 6, \overset{\uparrow}{8}, 9, \overset{\uparrow}{10}, 4$
 $\downarrow \quad \downarrow \quad \downarrow \quad \downarrow$
 $\overset{\circ}{i}_k, S_k \quad S_{k+1} = S_k + 1 \quad \overset{\circ}{i}_{k+1}, S_{k+1}$

Variables:

$$S_{k+1} = S_k + 1$$

$$\overset{\circ}{i}_{k+1} = \overset{\circ}{i}_{k-1}$$

Conditions:

$$(a[\overset{\circ}{i}_k] < 0) \Rightarrow \text{True}$$

$$(-a[\overset{\circ}{i}_k] == x) \Rightarrow \text{True}$$

$$\Rightarrow S_{k+1} = \sum_{j=\overset{\circ}{i}_{k+1}+1}^n (|a[\overset{\circ}{j}]| == x)$$

$$S_{k+1} = \sum_{j=\overset{\circ}{i}_{k-1}+1}^n (|a[\overset{\circ}{j}]| == x)$$

$$S_k = \sum_{j=\overset{\circ}{i}_k}^{n-1} (|a[\overset{\circ}{j}]| == x) - 1$$

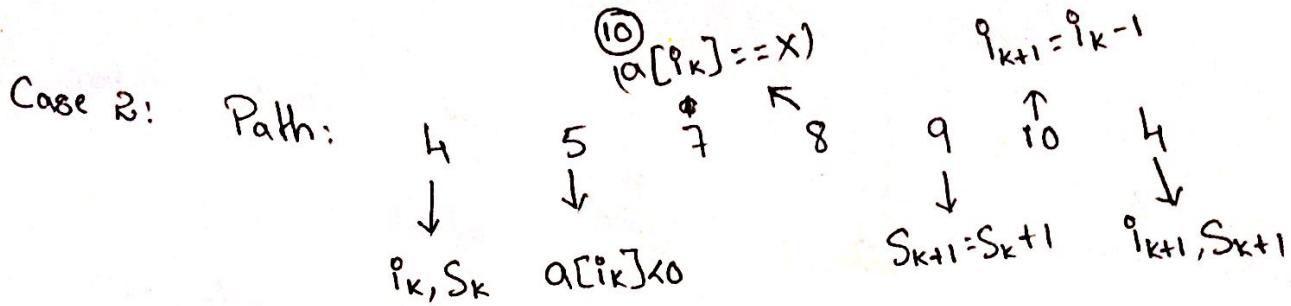
$$= \sum_{j=\overset{\circ}{i}_k}^{n-1} (|a[\overset{\circ}{j}]| == x) - \sum_{j=\overset{\circ}{i}_k}^{\overset{\circ}{i}_{k-1}} (a[\overset{\circ}{j}] == x)$$

Since Summation for $\overset{\circ}{i}_k$ is subtracted This is because

We can re-order the Summation to $j=\overset{\circ}{i}_{k-1}+1$

This can be regrouped to $(a[\overset{\circ}{i}_k] == x)$ is true.

$S_k = \sum_{j=\overset{\circ}{i}_{k+1}}^n (|a[\overset{\circ}{j}]| == x)$ is true based on
 the assumption in II



Variables:

$$S_{k+1} = S_k + 1$$

$$i_{k+1} = i_k - 1$$

Conditions:

$$a[i_k] < 0 \Rightarrow \text{False} \Rightarrow a[i_k] \geq 0$$

$$a[i_k] == x \Rightarrow \text{True.}$$

$$\Rightarrow S_{k+1} = \sum_{j=i_{k+1}}^n (|a[j]| == x)$$

$$S_{k+1} = \sum_{j=i_k-1+1}^n (|a[j]| == x)$$

$$S_k = \sum_{j=i_k}^{n-1} (|a[j]| == x) - 1$$

$$= \sum_{j=i_k}^{n-1} (|a[j]| == x) - \sum_{j=i_k}^{i_k-1} (a[j] == x)$$

↓

Since Summation for i_k is subtracted
we can re-ordered the Summation to

$$\sum_{j=i_{k+1}}$$

This is because $(a[i_k] == x)$
is true.

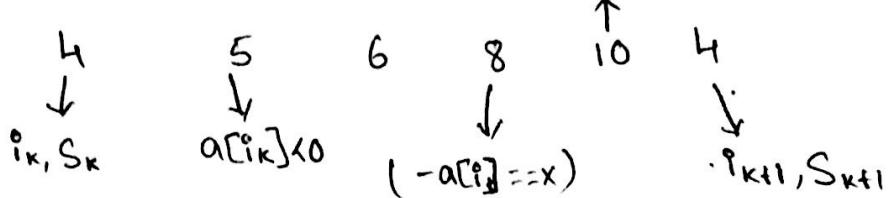
This can be regrouped to

$S_k = \sum_{j=i_{k+1}}^n (|a[j]| == x)$ is true based on the
assumption in II

(11)

$$i_{k+1} = i_k - 1$$

Case 3: Path:



Variables:

$$S_{k+1} = S_k$$

$$i_{k+1} = i_k - 1$$

Conditions:

$$a[i_k] < 0 \Rightarrow \text{True}$$

$$- a[i_k] == x \Rightarrow \text{False}$$

$$\Rightarrow S_{k+1} = \sum_{j=i_{k+1}+1}^n (|a[j]| == x)$$

$$S_k = \sum_{j=i_{k+1}-1}^n (|a[j]| == x)$$

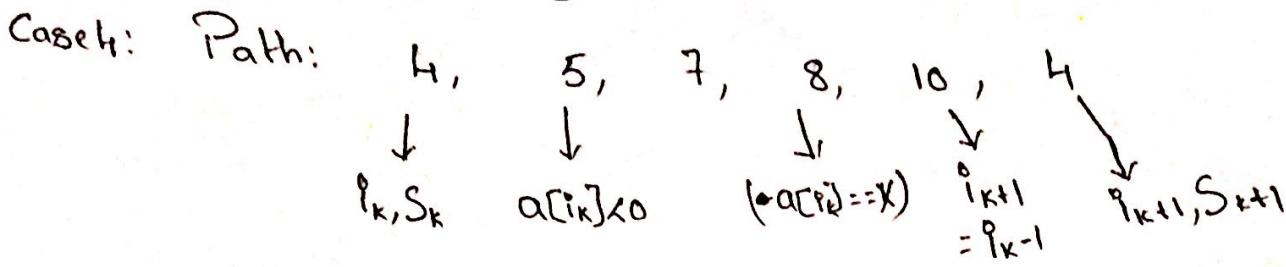
$$S_k = \sum_{j=i_{k+1}}^n (|a[j]| == x) + (a[i_k] == x)$$

$$S_k = \sum_{j=i_{k+1}}^n (|a[j]| == x)$$

Since $a[i_k] \neq x$ from Conditions
this returns zero

is true based on the
assumption in II.

(12)



Variables:

$$S_{k+1} = S_k.$$

$$i_{k+1} = i_k - 1$$

Conditions:

$$a[i_k] < 0 \Rightarrow \text{False}$$

$$a[i_k] == x \Rightarrow \text{False}$$

$$\Rightarrow S_{k+1} = \sum_{j=i_{k+1}+1}^n (|a[j]| == x)$$

$$S_k = \sum_{j=i_{k+1}-1}^{n-1} (|a[j]| == x)$$

$$S_k = \sum_{j=i_{k+1}}^n (|a[j]| == x) + (a[i_k] == x)$$

Since $a[i_k] \neq x$ from
Conditions this returns
zero

$$\Rightarrow S_k = \sum_{j=i_{k+1}}^n (|a[j]| == x)$$

is true based on

the assumption in II