# Auto Scaling of Key Value stores

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April 27, 2019

#### Outline

- Abstract
- 2 Introduction
- 3 Experiments
- 4 Current Picture
- 5 Further works

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#### Aim

- design a system to collect real time statistics on cluster based system.
- Collect statistics for different queries on different scaling configurations and try to find possible bottlenecks in the system.
- Providing an auto-scaling solution for key value stores.

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- Mostly In-memory extremy fast compared to traditional DBs.

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- Several load balancing schemes are used in horizontal scaling.

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- Why AutoScaling?
  - To prevent overprovisioning.

### Redis

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- Transaction, expiration time (BigTable).

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- After cluster meet, every node contain node hash slot mapping.

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- Moving hash slots through Cluster Bus.
- No down time with the help of MOVED Error.

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- Examples: Jedis, twemproxy, Redis Cluster.

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- Stats monitoring port.

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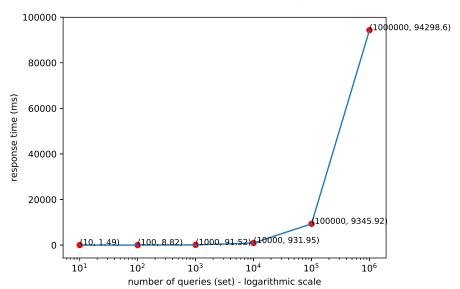
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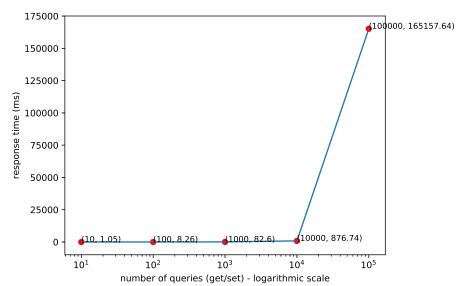
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- Resource limitation enforced using docker.

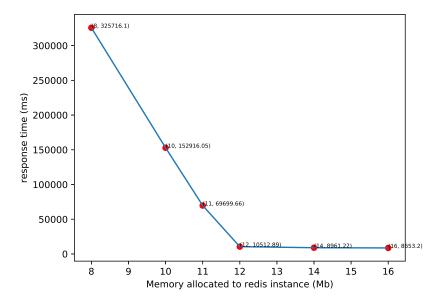
#### Single redis instance, with no memory/cpu restriction



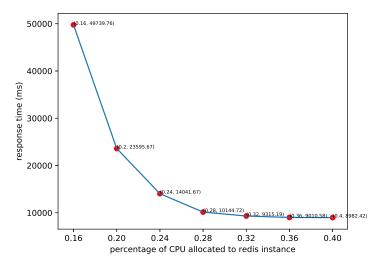
### Single redis instance, with max main memory = 10Mb



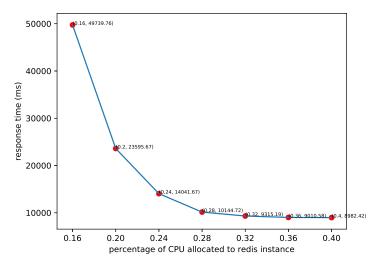
### 10^5 queries on redis instance, with varying main memory



#### 10^5 queries on redis instance, with varying CPU

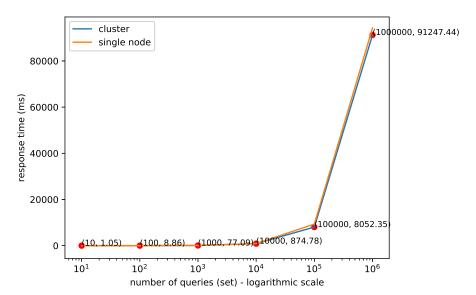


#### 10^5 queries on redis instance, with varying CPU

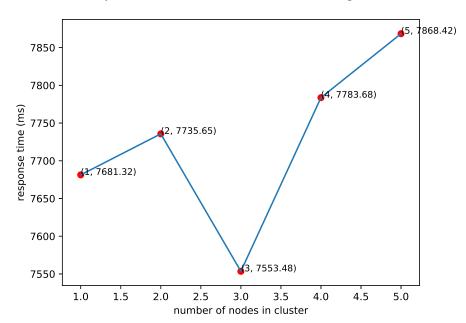


#### • fractional CPUs?

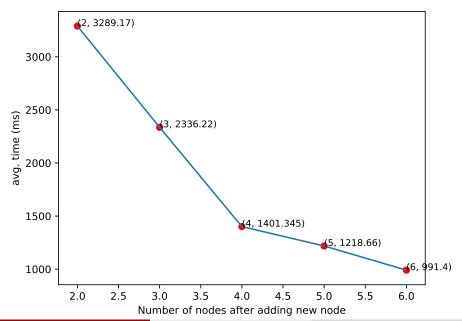
#### Single redis instance, compared with 6 Node cluster



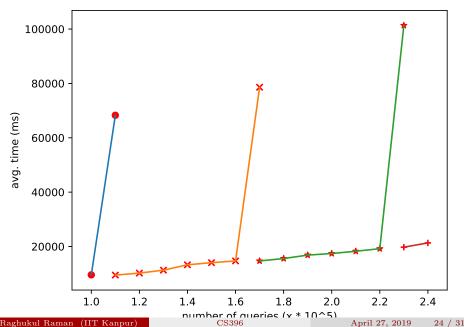
### 10^5 queries on different node configurations



### 10^5 keys present, setup time for adding new node



# Varying queries and number of nodes



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- Features used := startup time, resource utilization factor, etc.

Static metric based approach

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  - multiple requests to autoscale triggered by the same unschedulable pod may be invoked.

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- precise performance models can be automatically learned for distributed stream processing systems that can predict the execution performance of a job even before deployment.
- These models can be used to optimally schedule logically specified jobs onto available physical hardware.
- These models and the derived execution schedules can be refined online to dynamically adapt to unpredictable changes in the runtime environment or auto-scale with variations in job load.

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- Test some state of the art machine learning algorithms like Meta learning, Siamese nets etc.

# Thank You!