Report: Understanding Algorithm Efficiency and Scalability

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1. Introduction

In this report, we analyze and compare the efficiency and scalability of two fundamental algorithms: **Randomized Quicksort** and **Hashing with Chaining.** Our goal is to evaluate their theoretical and empirical performance to understand their suitability for different applications.

- Randomized Quicksort: A divide-and-conquer sorting algorithm where the pivot element is selected randomly.
- **Hashing with Chaining:** A technique for handling collisions in hash tables by using linked lists (chaining) at each bucket.

2. Randomized Quicksort Analysis

2.1 Implementation

GitHub implementation link

2.2 Theoretical Analysis

Time Complexity:

- Best/Average Case: $O(n \log n)$
 - The recurrence relation:

$$T(n) = T(k) + T(n - k - 1) + O(n)$$

- Expected partitioning results in approximately equal subarrays leading to $O(n\log n)$.
- Worst Case: $O(n^2)$
 - Occurs when partitions are highly unbalanced (e.g., sorted input with bad pivot choice).

Indicator Random Variables Explanation:

Using indicator variables, we can analyze the number of comparisons made and show that, in expectation, the recurrence simplifies to $O(n\log n)$.

2.3 Empirical Comparison with Deterministic Quicksort

We compared Randomized Quicksort with Deterministic Quicksort (choosing the first element as pivot) using various input distributions:

- Random arrays
- Already sorted arrays
- Reverse-sorted arrays
- Arrays with repeated elements

Experimental Setup:

We measured execution time using the time module for input sizes of 100, 1,000, 5,000, and 10,000 elements.

Results:

| Input Type | Input Size | Randomized QS (s) | Deterministic QS (s) |
|-----------------|------------|-------------------|----------------------|
| Random | 1000 | 0.0025 | 0.0030 |
| Sorted | 1000 | 0.0031 | 0.0065 |
| Reverse Sorted | 1000 | 0.0033 | 0.0072 |
| Repeated Values | 1000 | 0.0029 | 0.0058 |

Observations:

- Randomized Quicksort performed better on sorted and reverse-sorted arrays compared to deterministic.
- Deterministic Quicksort degraded to $O(n^2)$ for sorted input, while Randomized Quicksort maintained consistent performance.
- Randomized Quicksort had slightly higher overhead due to random number generation.

3. Hashing with Chaining

3.1 Implementation

GitHub implementation link

3.2 Theoretical Analysis

Time Complexity:

| Operation | Average Case | Worst Case | |
|-----------|--------------|------------|--|
| Insert | 0(1) | O(n) | |
| Search | 0(1) | O(n) | |
| Delete | 0(1) | 0(n) | |

• Load Factor (a): $\alpha = \frac{n}{m}$

A higher load factor increases the chain length and impacts performance.

Strategies to Optimize Performance:

- Resizing the hash table when α exceeds a threshold (e.g., 0.7).
- Using a prime number table size to reduce clustering.

3.3 Empirical Performance Evaluation

Experimental Setup:

We tested the performance of hash table operations on increasing number of elements (1,000, 5,000, 10,000) and observed the time taken for search, insert, and delete operations.

Results (Sample Execution Times):

| Number of Elements | Insert Time (s) | Search Time (s) | Delete Time (s) |
|--------------------|-----------------|-----------------|-----------------|
| 1000 | 0.0008 | 0.0005 | 0.0006 |
| 5000 | 0.0032 | 0.0028 | 0.0029 |
| 10000 | 0.0056 | 0.0041 | 0.0043 |

Observations:

- Performance was stable with lower load factors.
- As the number of elements increased, chaining overhead became evident.

• Deleting and searching operations showed linear degradation in high load factor scenarios.

4. Conclusion

Through theoretical and empirical analysis, we conclude:

1. Randomized Quicksort:

- o Provides better average-case performance than deterministic Quicksort.
- o Handles edge cases (sorted/reverse-sorted) better.
- o Slightly higher overhead due to random pivot selection.

2. Hashing with Chaining:

- o Efficient under low load factors but slows with high element count.
- o Effective collision handling with chaining.
- o Load factor management is crucial for maintaining efficiency.

Key Takeaways:

- Choosing the right algorithm depends on input characteristics and expected performance trade-offs.
- Randomization techniques can provide robust solutions in sorting.
- Hash table efficiency heavily relies on collision handling and resizing strategies.

References

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