



# **Digital Systems**

## **18B11EC213**

### **Module 1: Boolean Function Minimization Techniques and Combinational Circuits-7**

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# Standard Representations For Logical Functions

- The logical (Boolean) functions are expressed in terms of logical variables.
- Two standard forms for the representations of logical functions:

Sum-of-Products (SOP)

Product-of-Sums (POS)

- Literals: a variable on its own or in its complemented form. Examples:  $x$ ,  $x'$ ,  $y$ ,  $y'$
- Product Term: a single literal or a logical product (AND) of several literals.

Examples:  $x$ ,  $x.y.z'$ ,  $A'.B$ ,  $A.B$

## Cont..

- Sum Term: a single literal or a logical sum (OR) of several literals.

Examples:  $x$ ,  $x+y+z'$ ,  $A'+B$ ,  $A+B$

- Sum-of-Products (SOP) Expression: a product term or a logical sum (OR) of several product terms.

Examples:  $x$ ,  $x+y.z'$ ,  $x.y'+x'.y.z$ ,  $A.B+A'.B'$

- Product-of-Sums (POS) Expression: a sum term or a logical product (AND) of several sum terms.

Examples:  $x$ ,  $x.(y+z')$ ,  $(x+y').(x'+y+z)$ ,  $(A+B).(A'+B')$

## Cont..

- Every logical (Boolean) expression can either be expressed as SOP or POS expression.

Examples:

SOP:  $F1 = x'.y + x.y' + x.y.z$

POS:  $F2 = (x + y').(x' + y).(x' + z')$

# Minterm and Maxterm

- Consider two binary variables  $x$ ,  $y$ .
- Each variable may appear as itself or in complemented form as literals in a Boolean expression, i.e.,  $x$ ,  $x'$  and  $y$ ,  $y'$
- For two variables, there are four possible combinations with the AND operator, namely:

$$x'.y', x'.y, x.y', x.y$$

- These product terms are called the minterms.
- In general,  $n$  variables can give  $2^n$  minterms.

## Cont..

- In a similar fashion, a maxterm of  $n$  variables is the sum of  $n$  literals from the different variables.

Examples:  $x'+y'$ ,  $x'+y$ ,  $x+y'$ ,  $x+y$

- In general,  $n$  variables can give  $2^n$  maxterms.
- All the literals (or their complemented forms) should participate in the operations.

## Cont..

- The minterms and maxterms of two variables are denoted by m0 to m3 and M0 to M3 respectively.

Minterms				Maxterms	
x	y	term	notation	term	notation
0	0	$x'.y'$	m0	$x+y$	M0
0	1	$x'.y$	m1	$x+y'$	M1
1	0	$x.y'$	m2	$x'+y$	M2
1	1	$x.y$	m3	$x'+y'$	M3

- Each minterm is the complement of the corresponding maxterm and vice versa.

Example:  $m2 = x.y'$

$$m2' = (x.y')' = x' + (y')' = x'+y = M2$$

# Canonical Forms

- If each term (i.e., minterm and maxterm) in the SOP and POS forms/expressions contains all the literals (either in normal or complemented form), then these forms/expressions are known as the canonical SOP and canonical POS forms.

Example: For 3 variables a, b, c

Canonical SOP form:

$$Y = abc + abc' + ab'c' + a'bc$$

Canonical POS form:

$$Y = (a+b+c).(a+b'+c).(a+b+c').(a+b'+c')$$



# Cont..

- Canonical Form: Sum of Minterms

Example: Given a canonical SOP form

$$Y(a, b, c) = abc + abc' + ab'c' + a'bc$$

We can write Y as

$$\begin{aligned} Y &= m_7 + m_6 + m_4 + m_3 \\ &= m_3 + m_4 + m_6 + m_7 = \Sigma m(3, 4, 6, 7) \end{aligned}$$

where

$$a'bc = m_3$$

$$ab'c' = m_4$$

$$abc' = m_6$$

$$abc = m_7$$

# Cont..

Example: Obtain the expressions of F1, F2 and F3 from the following truth table.

x	y	z	F1	F2	F3
0	0	0	0	0	0
0	0	1	0	1	1
0	1	0	0	0	0
0	1	1	0	0	1
1	0	0	0	1	1
1	0	1	0	1	1
1	1	0	1	1	0
1	1	1	0	1	0

# Cont..

Obtain Sum-of-Minterms by summing the minterms of the function where the result is a 1.

$$F1 = x.y.z' = m6$$

$$F2 = x'.y'.z + x.y'.z' + x.y'.z + x.y.z' + x.y.z$$
$$= m1 + m4 + m5 + m6 + m7$$

$$= \Sigma m(1, 4, 5, 6, 7)$$

$$F3 = x'.y'.z + x'.y.z$$
$$+ x.y'.z' + x.y'.z$$
$$= m1 + m3 + m4 + m5$$
$$= \Sigma m(1, 3, 4, 5)$$

x	y	z	F1	F2	F3
0	0	0	0	0	0
0	0	1	0	1	1
0	1	0	0	0	0
0	1	1	0	0	1
1	0	0	0	1	1
1	0	1	0	1	1
1	1	0	1	1	0
1	1	1	0	1	0

# Cont..

- Canonical Form: Product of Maxterms
  - Maxterms are the sum terms.
  - Boolean functions can be expressed as Products-of-Maxterms.
  - For Boolean functions, the maxterms of a function are the terms for which the result is 0.

# Cont..

Example: Given a canonical POS form:

$$Y(a, b, c) = (a+b+c) (a+b+c') (a+b'+c) (a'+b+c')$$

We can write Y as

$$Y = M_0 \cdot M_1 \cdot M_2 \cdot M_5$$

(Since in maxterm, a is represented by 0 and a' by 1)

where  $(a+b+c) = M_0$

$$(a+b+c') = M_1$$

$$(a+b'+c) = M_2$$

$$(a'+b+c') = M_5$$

# Cont..

Example: Obtain the expressions of F2 and F3 from the following truth table.

x	y	z	F1	F2	F3
0	0	0	0	0	0
0	0	1	0	1	1
0	1	0	0	0	0
0	1	1	0	0	1
1	0	0	0	1	1
1	0	1	0	1	1
1	1	0	1	1	0
1	1	1	0	1	0

$$F2 = M_0 \cdot M_2 \cdot M_3 = \prod M(0, 2, 3) = (x+y+z).(x+y'+z).(x+y'+z')$$

$$\begin{aligned} F3 &= M_0 \cdot M_2 \cdot M_6 \cdot M_7 = \prod M(0, 2, 6, 7) \\ &= (x+y+z).(x+y'+z).(x'+y'+z).(x'+y'+z') \end{aligned}$$

# Cont..

- From the previous example (truth table), let us write the function F2 as the SOP form:

$$F2 = \sum m(1, 4, 5, 6, 7)$$

The complement function of F2 is:

$$\begin{aligned} F2' &= m0 + m2 + m3 \\ &= \sum m(0, 2, 3) \end{aligned}$$

x	y	z	F2	F2'
0	0	0	0	1
0	0	1	1	0
0	1	0	0	1
0	1	1	0	1
1	0	0	1	0
1	0	1	1	0
1	1	0	1	0
1	1	1	1	0

- This means that the functions F2 and F2' jointly contain all possible minterms.

# Cont..

From the previous slide, we obtained

$$F2' = m_0 + m_2 + m_3$$

Therefore,

$$F2 = (m_0 + m_2 + m_3)'$$

$$= m_0' \cdot m_2' \cdot m_3'$$

DeMorgan's theorem

$$= M_0 \cdot M_2 \cdot M_3$$

Since  $m_x' = M_x$

$$= \prod M(0, 2, 3)$$

$$\Rightarrow F2 = \sum m(1, 4, 5, 6, 7) = \prod M(0, 2, 3)$$

□ Every Boolean function can be expressed as either Sum-of-Minterms or Product-of-Maxterms.



# Conversion of Canonical Forms

- Sum-of-Minterms to Product-of-Maxterms

Example: Given  $F1(A, B, C) = \sum m(3, 4, 5, 6, 7)$  **SOP**

then  $F1(A, B, C) = \prod M(0, 1, 2)$  **POS**

(Replace minterm indices with indices not already used)

- Product-of-Maxterms to Sum-of-Minterms

Example: Given  $F2(A, B, C) = \prod M(0, 3, 5, 6)$  **POS**

then  $F2(A, B, C) = \sum m(1, 2, 4, 7)$  **SOP**

(Replace maxterm indices with indices not already used)

## Cont..

- Sum-of-Minterms of  $F \Rightarrow$  Product-of-Maxterms of  $F'$

Example: Given  $F1(A, B, C) = \sum m(3, 4, 5, 6, 7)$

then  $F1'(A, B, C) = \prod M(3, 4, 5, 6, 7)$

(use the same indices as in  $F1$ )

- Product-of-Maxterms of  $F \Rightarrow$  Sum-of-Minterms of  $F'$

Example: Given  $F1(A, B, C) = \prod M(0, 1, 2)$

then  $F1'(A, B, C) = \sum m(0, 1, 2)$

(use the same indices as in  $F1$ )

## Cont..

- Sum-of-Minterms of  $F \Rightarrow$  Sum-of-Minterms of  $F'$

Example: Given  $F1(A, B, C) = \sum m(3, 4, 5, 6, 7)$

then  $F1'(A, B, C) = \sum m(0, 1, 2)$

(list the indices not already used in  $F$ )

- Product-of-Maxterms of  $F \Rightarrow$  Product-of-Maxterms of  $F'$

Example: Given  $F1(A, B, C) = \prod M(0, 1, 2)$

then  $F1'(A, B, C) = \prod M(3, 4, 5, 6, 7)$

(list the indices not already used in  $F$ )

# References

- M. M. Mano, *Digital Logic and Computer Design*, 5th ed., Pearson Prentice Hall, 2013.
- R. P. Jain, *Modern Digital Electronics*, 4th ed., Tata McGraw-Hill Education, 2009.