

## **Tutorial Sheet – Odd Semester 2022**

### **15B11CI212 – Theoretical Foundations of Computer Science**

#### **Tutorial ... 12**

#### **Automata Theory**

##### **1. DFAs**

Draw Deterministic Finite Automata to accept the following sets of strings over the alphabet  $\{0,1\}$ :

- a. All strings that contain exactly 4 "0"s.
- b. All strings ending in "1101".
- c. All strings containing exactly 4 "0"s and at least 2 "1"s.
- d. All strings whose binary interpretation is divisible by 5.
- e. All strings that contain the substring 0101.
- f. All strings that start with 0 and have odd length or start with 1 and have even length.
- g. All strings that don't contain the substring 110.
- h. All strings of length at most five.
- i. All strings where every odd position is a 1.

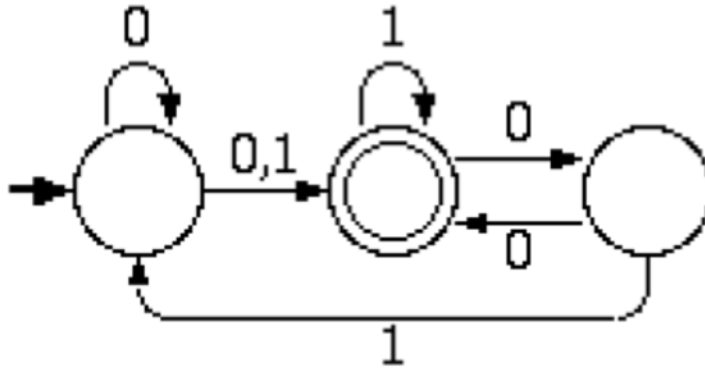
##### **2. NFAs**

Draw Non-deterministic Finite Automata with the specified number of states to accept the following sets:

- a. All strings containing exactly 4 "0"s or an even number of "1"s. (8 states)
- b. All strings such that the third symbol from the right end is a "0". (4 states)
- c. All strings such that some two zeros are separated by a string whose length is  $4i$  for some  $i \geq 0$ . (6 states)
- d. All strings that contain the substring 0101. (5 states)
- e. All strings that contain an even number of zeros or exactly two ones. (6 states)
- f. The language  $0^*1^*0^*0$ . (3 states)

##### **3. Converting NFAs to DFAs**

Consider the following NFA over the alphabet  $\{0,1\}$ :



- a. Convert this NFA to a minimal DFA.
- b. Write a regular expression for the set the machine accepts.

#### 4. Discrete Math Review – Proofs

Analyze the two languages below. They are two descriptions of the same language – strings of balanced parentheses.

Language 1: The set of strings where each string  $w$  has an equal number of zeros and ones; and any prefix of  $w$  has at least as many zeros as ones.

Language 2: The set of strings defined inductively as follows: if  $w$  is in the set then  $0w1$  is also in the set; if  $u$  and  $v$  are in the set then so is  $uv$ ; and the empty string is in the set.

- a. Prove that every string in Language 2 is contained in Language 1.
- b. Extra Credit: Prove they are equal (i.e. Language 1 is also contained in Language 2).

#### 5. Closure Problems

You may use examples to illustrate your proofs.

- a. Prove that if  $L_1$  is regular and  $L_2$  is regular then so is  $L_1 - L_2$  (the set of all strings in  $L_1$  but not in  $L_2$ ).
- b. Prove that if  $L$  is regular then  $\text{Prefix}(L)$  is regular.  $\text{Prefix}(L)$  is the set of all strings which are a proper prefix of a string in  $L$ .
- c. Prove that Regular Sets are closed under MIN.  $\text{MIN}(R)$ , where  $R$  is a regular set, is the set of all strings  $w$  in  $R$  where every proper prefix of  $w$  is

in not in R. (Note that this is not simply the complement of PREFIX).

d. Prove that Regular Sets are NOT closed under infinite union. (A counterexample suffices).

e. What about infinite intersection?

## 6. Regular Expressions

Write regular expressions for each of the following languages over the alphabet  $\{0,1\}$ . Provide justification that your regular expression is correct.

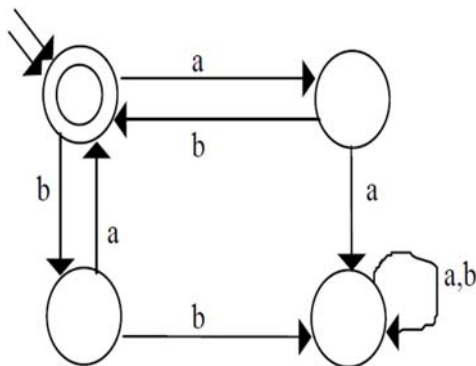
a. The set of all strings in which every pair of adjacent zeros appears before any pair of adjacent ones.

b. The set of all strings not containing 101 as a substring.

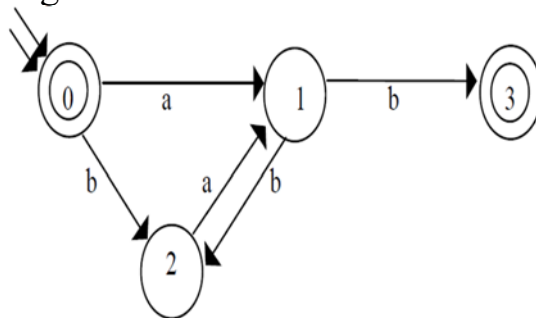
c. The set of all strings with at most one pair of consecutive zeros and one pair of consecutive ones.

## 7. Converting Finite Automata to Regular Expressions

a. Write a regular expression for the language recognized by the following FSM:



b. Consider the following FSM M:



- (i) Write a regular expression for the language accepted by M.
- (ii) Give a deterministic FSM that accepts the complement of the language accepted by M.

## 8. Regular Expression Identities

Prove (give at least a few words of justification), or disprove (by counterexample) that each pair of regular expressions represent the same language. Assume that  $r$ ,  $s$  and  $t$  represent regular expressions over the alphabet  $\{0,1\}$ .

- a.  $r(s + t)$  and  $rs + rt$
- b.  $(r^*)^*$  and  $r^*$
- c.  $(r + s)^*$  and  $r^*s^*$

## 9. Final States

- a. Explain why every NFA can be converted to an equivalent one that has a single final state.
- b. Give a counterexample to show that this is not true for DFA's.
- c. Extra Credit: Describe the languages that are generated from a DFA with just one final state.

## 10. Problems

- a. Draw a Finite Automaton to accept the following regular expression and succinctly describe the set in English.

$$[00 + 11 + (01 + 10)(00 + 11)^*(01 + 10)]^*$$

## 11. More Machines

Draw a finite state machine that accepts the complement of the language accepted by the non-deterministic machine below:

