

Assignment 2 - Emirps

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COMP2111 18s1

1 Task 1 - Specification Statement

Notes,x:

- Write neatly
- make sure grammar is correct
- look at examples for default spec structure.

Define an Emirp using 2 functions - reverse and prime. Make these functions match with their given specs in order to help prove implications.

Pre condition,x: n is a positive number - $n > 0$

Post condition $EMIRP(r, n)$ where r is the n^{th} emirp (where emirp is as defined above).
Therefore our program can be specified by,x:

proc EMIRP(**value** n , **result** r) ·
 $\sqcup n, r, x : [n > 0, Emirp(r, n)] \neg(1)$

2 Task 2 - Derivation

$$\begin{aligned}
& \text{proc EMIRP}(\text{value } n, \text{result } r) \cdot \sqcup n, r, x : [n > 0, \text{Emirp}(r, n)] \quad \textcolor{red}{\dashv(1)} \\
(1) \sqsubseteq & \quad \langle \text{c-frame} \rangle \\
& \sqcup r, x : [n > 0, \text{Emirp}(r, n)] \quad \textcolor{red}{\dashv(2)} \\
(2) \sqsubseteq & \quad \langle \text{i-loc} \rangle \\
& \sqcup i, r, x : [n > 0, \text{Emirp}(r, n)] \quad \textcolor{red}{\dashv(3)} \\
(3) \sqsubseteq & \quad \langle \text{seq} \rangle \\
& \sqcup i, x, r : [n > 0, i = 1 \wedge x = 13 \wedge n > 0] \quad \textcolor{red}{\dashv(4)}; \\
& \sqcup i, x : [i = 1 \wedge x = 13 \wedge n > 0, \text{Emirp}(r, n)] \quad \textcolor{red}{\dashv(5)} \\
(4) \sqsubseteq & \quad \langle \text{c-frame} \rangle \\
& i, x : [n > 0, i = 1 \wedge x = 13 \wedge n > 0] \\
& \sqsubseteq \quad \langle \text{ass - (1)} \rangle \\
& i := 1 \\
& x := 13 \\
(5) \sqsubseteq & \quad \langle \text{seq} \rangle \\
& \sqcup i, r, x : [i = 1 \wedge n > 0, \text{Inv}] \quad \textcolor{red}{\dashv(6)}; \\
& \sqcup i, r, x : [\text{Inv}, \text{Inv} \wedge i = n] \quad \textcolor{red}{\dashv(7)}; \\
& \sqcup i, r, x : [\text{Inv} \wedge i = n, \text{Emirp}(r, n)] \quad \textcolor{red}{\dashv(8)} \\
(6) \sqsubseteq & \quad \langle \text{w-pre, c-frame - (2)} \rangle \\
& r, x : [\text{Inv}^{[13/x]}, \text{Inv}] \\
& \sqsubseteq \quad \langle \text{ass - (3)} \rangle \\
& r := 13 \\
(7) \sqsubseteq & \quad \langle \text{while} \rangle \\
& \text{while } i \neq n \text{ do} \\
& \quad \sqcup i, r, x : [\text{Inv} \wedge i \neq n, \text{Inv}] \quad \textcolor{red}{\dashv(9)} \\
& \text{od;} \\
(8) \sqsubseteq & \quad \langle \text{w-pre, c-frame - (4)} \rangle \\
& \sqcup r, x : [\text{EMIRP}(r, n), \text{EMIRP}(r, n)] \quad \textcolor{red}{\dashv(10)} \\
& \sqsubseteq \quad \langle \text{skip - (5)} \rangle \\
& \text{skip} \\
(9) \sqsubseteq & \quad \langle \text{seq} \rangle \\
& \sqcup r, x : [\text{Inv} \wedge i \neq n, \text{Inv}^{[x+1/x]}] \quad \textcolor{red}{\dashv(10)}; \\
& \sqcup r, x : [\text{Inv}^{[x+1/x]}, \text{Inv}] \quad \textcolor{red}{\dashv(11)}
\end{aligned}$$

$$\begin{aligned}
(10) &\sqsubseteq \langle \text{ass - (6)} \rangle \\
&x := x + 1 \\
(11) &\sqsubseteq \langle \text{i-loc, seq} \rangle \\
&\sqsubseteq a, i, r, x : [\text{Inv}^{[x+1/x]}, \text{Inv}^{[x+1/x]} \wedge a = 1] \neg(12); \\
&\sqsubseteq a, i, r, x : [\text{Inv}^{[x+1/x]} \wedge a = 1, \text{Inv}] \neg(13) \\
(12) &\sqsubseteq \langle \text{c-frame} \rangle \\
&a, x : [\text{Inv}^{[x+1/x]}, \text{Inv}^{[x+1/x]} \wedge a = 1] \\
&\sqsubseteq \langle \text{ass - (7)} \rangle \\
&a := 1 \\
(13) &\sqsubseteq \langle \text{seq} \rangle \\
&\sqsubseteq a, i, r, x : \left[\begin{array}{l} \text{Inv}^{[x+1/x]} \wedge a = 1, (a = 1 \wedge \neg \exists k \in 2..(x-1) (x \bmod k = 0)) \\ \vee (a = 0 \wedge \exists k \in 2..(x-1) (x \bmod k = 0)) \end{array} \right] \neg(14); \\
&\sqsubseteq a, i, r, x : \left[\begin{array}{l} (a = 1 \wedge \neg \exists k \in 2..(x-1) (x \bmod k = 0)) \\ \vee (a = 0 \wedge \exists k \in 2..(x-1) (x \bmod k = 0)), \text{Inv} \end{array} \right] \neg(15) \\
(14) &\sqsubseteq \langle \text{w-pre - (8)} \rangle \\
&a, x : [a > 1 \wedge x > 0, \text{post}(14)] \\
&\sqsubseteq \langle \text{proc} \rangle \\
&\text{isPrime}(x, a) \\
(15) &\sqsubseteq \langle \text{if} \rangle \\
&\text{if } a = 1 \\
&\quad \text{then } \sqsubseteq a, i, r, x : [a = 1 \wedge \text{pre}(15), \text{post}(15)] \neg(16) \\
&\quad \text{else } \sqsubseteq p, x : [a \neq 1 \wedge \text{pre}(15), \text{post}(15)] \neg(17) \\
&\quad \text{fi} \\
(16) &\sqsubseteq \langle \text{i-loc} \rangle \\
&a, i, r, s, x : [\text{pre}(16), \text{post}(16)] \\
&\sqsubseteq \langle \text{seq} \rangle \\
&\sqsubseteq s, x : [\text{pre}(16), s = 0] \neg(18); \\
&\sqsubseteq a, i, r, s, x : [s = 0, \text{post}(16)] \neg(19) \\
(17) &\sqsubseteq \langle \text{c-frame, w-pre - (9)} \rangle \\
&i, r : [\text{Inv}, \text{Inv}] \\
&\sqsubseteq \langle \text{ass - (10)} \rangle \\
&\text{skip} \\
(18) &\sqsubseteq \langle \text{ass - (11)} \rangle \\
&s := 0 \\
(19) &\sqsubseteq \langle \text{seq} \rangle \\
&\sqsubseteq s, x : [\text{pre}(19), \text{reversen function post condition}] \neg(20); \\
&\sqsubseteq a, i, r, s, x : [\text{reversen function post condition}, \text{post}(19)] \neg(21)
\end{aligned}$$

$$\begin{aligned}
(20) &\sqsubseteq \langle \text{w-pre - (12)} \rangle \\
&\quad s, x : [\text{reverse function pre condition, reversen function post condition}] \\
&\sqsubseteq \langle \text{proc} \rangle \\
&\quad \text{reversen}(x, s) \\
(21) &\sqsubseteq \langle \text{i-loc, seq} \rangle \\
&\quad \sqsubseteq a, i, r, s, b, x : [\text{pre}(21), \text{pre}(21) \wedge b = 1] \neg(22); \\
&\quad \sqsubseteq a, i, r, s, b, x : [\text{pre}(21) \wedge b = 1, \text{post}(21)] \neg(23) \\
(22) &\sqsubseteq \langle \text{c-frame, ass - (13)} \rangle \\
&\quad b := 1 \\
(23) &\sqsubseteq \langle \text{seq} \rangle \\
&\quad \sqsubseteq a, i, r, s, b, x : \left[\begin{array}{l} \text{pre}(21) \wedge b = 1, (a = 1 \wedge \neg \exists k \in 2..(x-1) (x \bmod k = 0)) \\ \vee (a = 0 \wedge \exists k \in 2..(x-1) (x \bmod k = 0)) \end{array} \right] \neg(24); \\
&\quad \sqsubseteq a, i, r, s, b, x : \left[\begin{array}{l} (a = 1 \wedge \neg \exists k \in 2..(x-1) (x \bmod k = 0)) \\ \vee (a = 0 \wedge \exists k \in 2..(x-1) (x \bmod k = 0)), \text{post}(21) \end{array} \right] \neg(25) \\
(24) &\sqsubseteq \langle \text{w-pre - (14)} \rangle \\
&\quad a, i, r, s, b, x : \left[\begin{array}{l} s > 0 \wedge b = 1, (a = 1 \wedge \neg \exists k \in 2..(x-1) (x \bmod k = 0)) \\ \vee (a = 0 \wedge \exists k \in 2..(x-1) (x \bmod k = 0)) \end{array} \right] \\
&\sqsubseteq \langle \text{proc} \rangle \\
&\quad \text{isPrime}(x, a) \\
(25) &\sqsubseteq \langle \text{if} \rangle \\
&\quad \text{if } b = 1 \wedge s \neq x \\
&\quad \text{then } \sqsubseteq i, x : [b = 1 \wedge s \neq x \wedge \text{pre}(25), \text{post}(25)] \neg(26) \\
&\quad \text{else } \sqsubseteq i, r, a, s, b, x : [(b \neq 1 \vee s = x) \wedge \text{pre}(25), \text{post}(25)] \neg(27) \\
&\quad \text{fi} \\
(26) &\sqsubseteq \langle \text{c-frame, w-pre- (15)} \rangle \\
&\quad a, i, r, s, b, x : [\text{Inv}^{i+1}/i][x/r], \text{Inv}] \\
&\sqsubseteq \langle \text{ass - (16)} \rangle \\
&\quad i := i + 1 \\
&\quad r := r \\
(27) &\sqsubseteq \langle \text{c-frame, w-pre- (17)} \rangle \\
&\quad a, i, r, s, b, x : [\text{Inv}, \text{Inv}] \\
&\sqsubseteq \langle \text{ass - (18)} \rangle \\
&\quad \text{skip}
\end{aligned}$$

We gather the code for the procedure body of EMIRP,x:

We have derived our code. However we need to prove **some** refinements.

2.1 Implication 1: $(4) \sqsubseteq i := 1$

$$\begin{aligned}
& K \wedge i \geq 1 \wedge a[i][l] = a[i-1][l] \wedge j = 1 \wedge a[i][l] = 0 \\
\Rightarrow & \langle K \wedge B \wedge C \wedge D \wedge E \Rightarrow K \wedge C \wedge E. \text{Expanding } K \rangle \\
& I \wedge j \in (0, 1) \wedge l \in 0..y, a[i][y] = 0 \wedge \forall w \in 0..(l-1) (a[i][w] = a[i-1][w]) \wedge \\
& a[i][l] = a[i-1][l] \wedge a[i][l] = 0 \\
\Rightarrow & \langle \text{since true for } w \text{ in } 0..l-1 \text{ and also for } l \rangle \\
& I \wedge j \in (0, 1) \wedge l \in 0..y, a[i][y] = 0 \wedge \forall w \in 0..l (a[i][w] = a[i-1][w] \wedge a[i][l] = 0) \\
\Rightarrow & \langle \text{separating for } l \text{ in } 0..y-1 \text{ and } l=y \rangle \\
& I \wedge j \in (0, 1) \wedge l \in 0..y-1, a[i][y] = 0 \wedge \forall w \in 0..l (a[i][w] = a[i-1][w]) \wedge a[i][l] = 0 \vee \\
& I \wedge j \in (0, 1) \wedge l = y \wedge a[i][l] = 0, a[i][y] = 0 \wedge \forall w \in 0..l (a[i][w] = a[i-1][w]) \\
\Rightarrow & \langle \text{since } a[i][l] = a[i][y] = 0 \text{ when } l=y \rangle \\
& I \wedge j \in (0, 1) \wedge l \in 0..y-1, a[i][y] = 0 \wedge \forall w \in 0..l (a[i][w] = a[i-1][w]) \vee TRUE \\
\Leftrightarrow & \langle \text{definitions of } I \text{ and substitution, } l \text{ in } 0..y-1 \text{ means } (l+1) \text{ in } 0..y \rangle \\
& K^{[l+1/l][0/j]}
\end{aligned}$$

3 Task 3 - C Code

```

1  #include <stdio.h>
2  #include "reverse.h"
3
4  unsigned long emirp(unsigned long n);
5  void isPrime(unsigned long r, int *a);
6
7  int main (int argc, char* argv[]){
8      unsigned long n;
9      if(scanf("%lu", &n)==1)
10         printf("%lu\n",emirp(n));
11 }
12
13 /*
14  var i := 1
15  r := 13
16  while i != n do
17     r := r + 1
18     var a := 1

```

```

19      isPrime(r,a)
20      if a = 1 then
21          var s := 0
22          reversen(r,s)
23          var b := 1
24          isPrime(s, b)
25          if b = 1 && s != r then
26              i = i + 1
27      od
28
29  */
30  unsigned long emirp(unsigned long n) {
31      int i = 1;
32      unsigned long r = 13;
33      while (i != n) {
34          r = r + 1;
35          int a = 1;
36          isPrime(r, &a);
37          if (a == 1) {
38              unsigned long s = 0;
39              reversen(r, &s);
40              int b = 1;
41              isPrime(s, &b);
42              if (b == 1 && s != r) {
43                  i = i + 1;
44              }
45          }
46      }
47      return r;
48  }
49
50  /*
51  var j := 0
52  while j != r do
53      if r mod j = 0 then
54          a = 0
55      j := j + 1
56  od
57  */
58  void isPrime(unsigned long r, int *a) {
59      unsigned long j = 2;
60      while (j != r) {
61          if (r % j == 0) {
62              *a = 0;

```

```
63     }  
64     j = j + 1;  
65 }  
66 }
```

- Write something about how the C code relates.
- Compare with examples