
Compiler Design Project

This project creates a **Graphical User Interface (GUI) application** for analyzing and visualizing different components of a compiler's functionality using Python and the Tkinter library.

LEXER :This code is a **lexer** that converts source code into tokens for a compiler or interpreter. It:

```
1  TOKENS = [  
2      ("ASSIGN", "="),  
3      ("PLUS", "+"),  
4      ("MINUS", "-"),  
5      ("MULT", "*"),  
6      ("DIV", "/"),  
7      ("LARGER THAN", ">"),  
8      ("LESSTHAN", "<"),  
9      ("EQUALS", "=="),  
10     ("NOTEQUALS", "!="),  
11     ("LPAREN", "("),  
12     ("RPAREN", ")"),  
13     ("LBRACE", "{"),  
14     ("RBRACE", "}"),  
15     ("LBRACKET", "["),  
16     ("RBRACKET", "]"),  
17     ("COLON", ":"),  
18     ("COMMA", ","),  
19     ("NUMBER", "0123456789"),  
20     ("WHITESPACE", " \t\n"),  
21 ]  
22  
23 KEYWORDS = {"if", "else", "for", "print", "in", "range"}  
24  
25 def is_identifier_start(char):  
26     return char.isalpha() or char == "_"  
27  
28 def is_identifier_part(char):  
29     return char.isalnum() or char == "_"
```

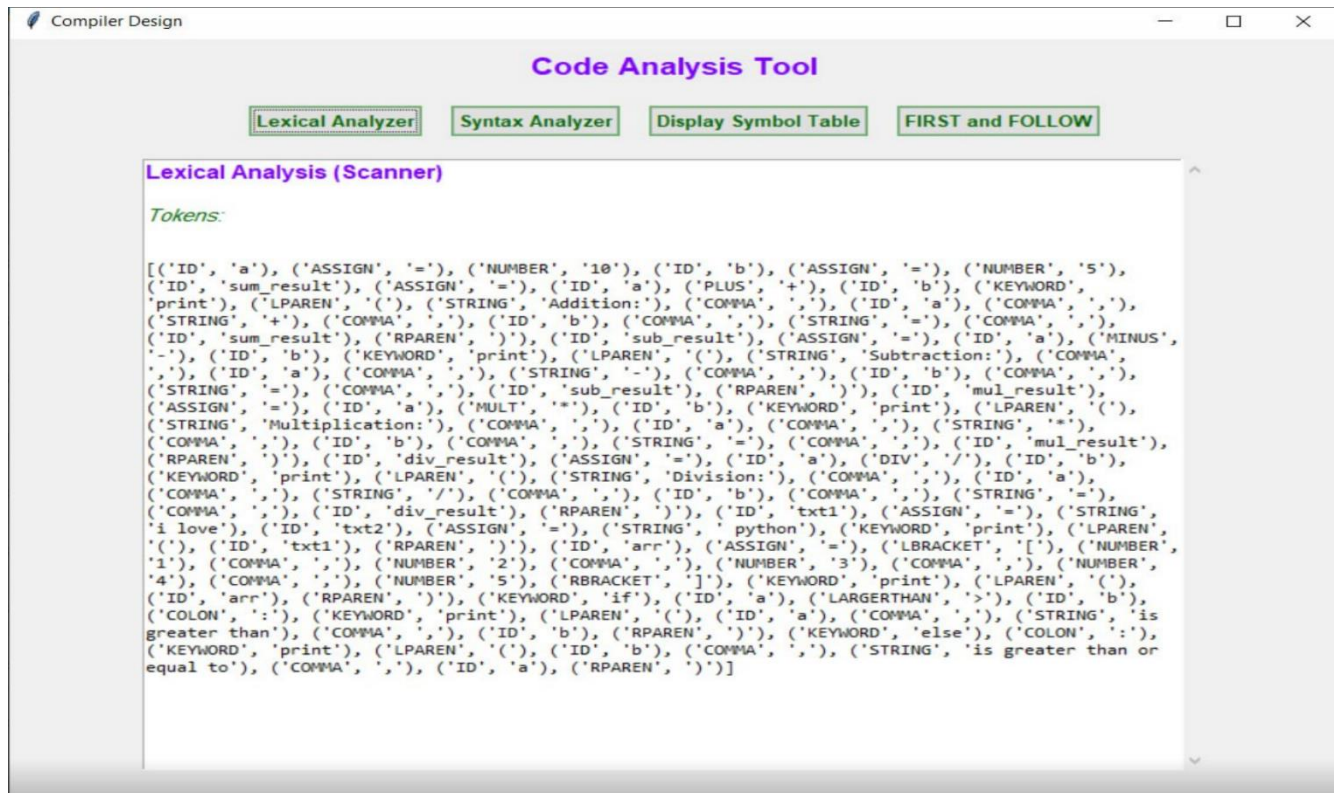
```
31 def tokenize(code):
32     tokens = []
33     i = 0
34     line = 1
35     col = 1
36
37     while i < len(code):
38         char = code[i]
39
40         if char in " \t\n":
41             if char == "\n":
42                 line += 1
43                 col = 1
44             else:
45                 col += 1
46             i += 1
47             continue
48
49         if is_identifier_start(char):
50             identifier = char
51             i += 1
52             col += 1
53             while i < len(code) and is_identifier_part(code[i]):
54                 identifier += code[i]
55                 i += 1
56                 col += 1
57             if identifier in KEYWORDS:
58                 tokens.append(("KEYWORD", identifier))
59             else:
60                 tokens.append(("ID", identifier))
61             continue
```

```

63     # quote_type = ' ' \
64     if char in '"':
65         quote_type = char
66         string_literal = ""
67         i += 1
68         col += 1
69         while i < len(code) and code[i] != quote_type:
70             # مافيش قافله استرينج
71             if code[i] == "\n":
72                 raise ValueError(f"Unterminated string literal at line {line}, col {col}")
73             string_literal += code[i]
74             i += 1
75             col += 1
76         if i >= len(code) or code[i] != quote_type:
77             raise ValueError(f"Unterminated string literal at line {line}, col {col}")
78         i += 1
79         col += 1
80         tokens.append(("STRING", string_literal))
81         continue
82
83     #multi-character operators: == , !=
84     if i + 1 < len(code) and code[i:i + 2] in {"==", "!="}:
85         tokens.append(("OPERATOR", code[i:i + 2]))
86         i += 2
87         col += 2
88         continue
89
90     # single-character tokens
91     for type, character in TOKENS:
92         if char in character:
93             if type == "NUMBER":
94                 number = char
95                 i += 1
96                 col += 1
97                 while i < len(code) and code[i] in "0123456789":
98                     number += code[i]
99                     i += 1
100                    col += 1
101                    tokens.append((type, number))
102             else:
103                 tokens.append((type, char))
104                 i += 1
105                 col += 1
106             break
107         else:
108             raise ValueError(f"Unexpected character '{char}' at line {line}, col {col}")
109
110     return tokens
111

```

The result is a list of tokens for further processing in the compilation or interpretation process.



Defines Tokens: Recognizes operators (+, =), symbols ({, }), numbers, strings, and keywords (if, else).

Processes Code:

- Skips whitespace.

- Identifies keywords or variable names (identifiers).

- Detects numbers and string literals.

- Handles multi-character (==, !=) and single-character tokens.

Error Handling: Raises an error for unexpected characters or unterminated strings.

PARSER : This code defines a parser that takes a list of tokens (produced by a lexer) and constructs a parse tree representing the structure of a program.

Here's a brief explanation:

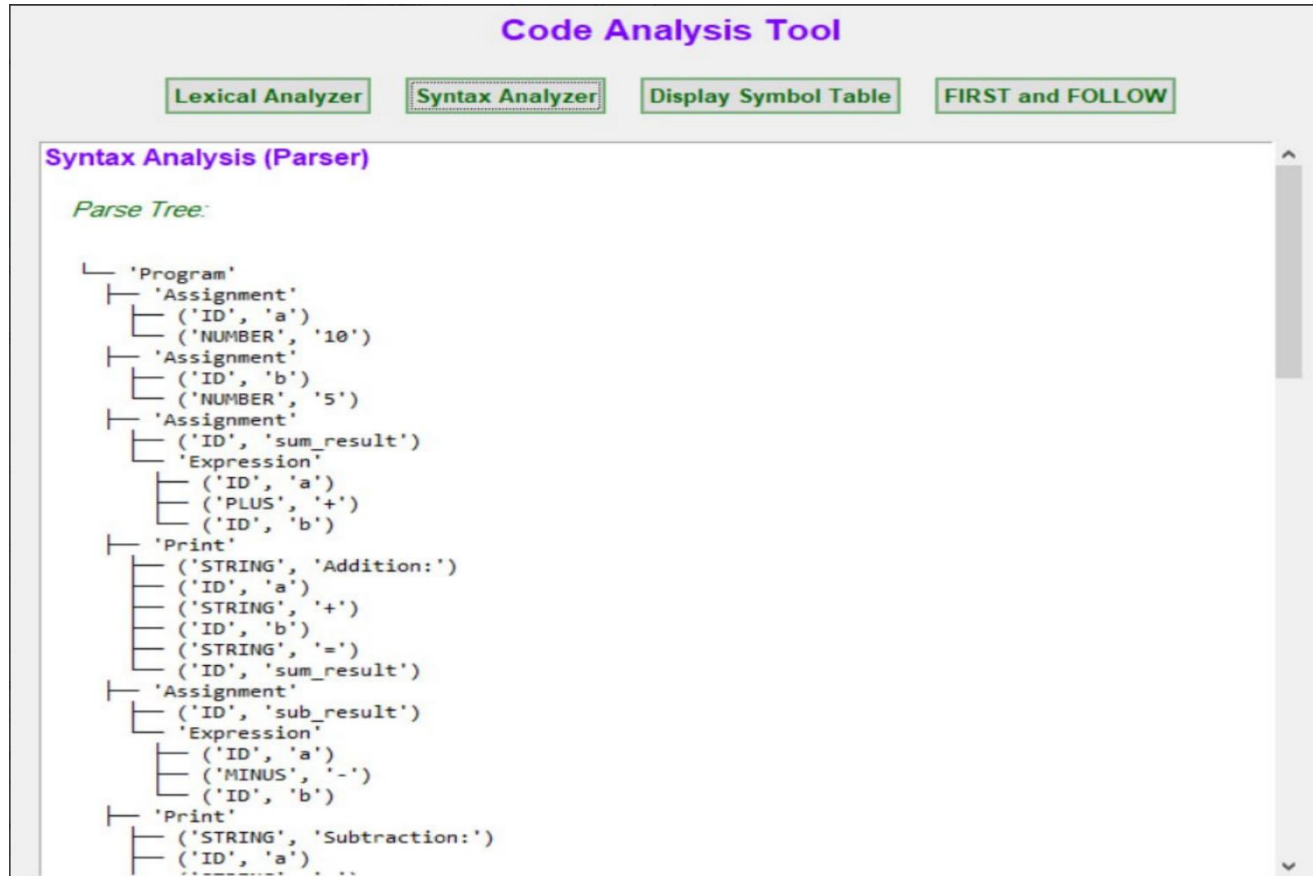
```
1 class ParseNode:
2     def __init__(self, value, children=None):
3         self.value = value
4         self.children = children or []
5
6     def __str__(self, level=0, is_last=True):
7         edge = "└─" if is_last else "├─"
8         ret = "  " * level + edge + repr(self.value) + "\n"
9
10        for i, child in enumerate(self.children):
11            ret += child.__str__(level + 1, is_last=(i == len(self.children) - 1))
12
13        return ret
14
15
16 class Parser:
17     def __init__(self, tokens):
18         self.tokens = tokens
19         self.pos = 0
20
21     def current_token(self):
22         return self.tokens[self.pos] if self.pos < len(self.tokens) else None
23
24     def match(self, expected_type):
25         token = self.current_token()
26         if token and token[0] == expected_type:
27             self.pos += 1
28             return token
29         else:
30             raise SyntaxError(f"Expected {expected_type}, found {token}")
31
32     def parse_program(self):
33         return ParseNode("Program", self.parse_statements())
34
35     def parse_statements(self):
36         statements = []
37         while self.current_token():
38             statements.append(self.parse_statement())
39         return statements
40
41     def parse_statement(self):
42         token = self.current_token()
43         if token[0] == "ID":
44             return self.parse_assignment()
45         elif token[0] == "KEYWORD" and token[1] == "if":
46             return self.parse_if_statement()
47         elif token[0] == "KEYWORD" and token[1] == "else":
48             return self.parse_else_statement()
49         elif token[0] == "KEYWORD" and token[1] == "print":
50             return self.parse_print_statement()
51         else:
52             raise SyntaxError(f"Unexpected statement: {token}")
53
54     def parse_assignment(self):
55         id_token = self.match("ID")
56         self.match("ASSIGN")
57         expr = self.parse_expression()
58         return ParseNode("Assignment", [ParseNode(id_token), expr])
59
60     def parse_else_statement(self):
61         self.match("KEYWORD")
62         self.match("COLON")
63         else_body = self.parse_statements()
64         return ParseNode("Else", else_body)
65
66     def parse_if_statement(self):
67         self.match("KEYWORD")
68         condition = self.parse_expression()
69         self.match("COLON")
70         body = self.parse_statements()
71
72         else_node = None
73         if self.current_token() and self.current_token()[0] == "KEYWORD" and self.current_token()[1] == "else":
74             else_node = self.parse_else_statement()
75
76         children = [
77             condition,
78             ParseNode("Body", body)
79         ]
80
81         if else_node:
82             children.append(ParseNode("Else", [else_node])) # Only add the Else node if it exists
83
84         return ParseNode("IfStatement", children)
```

```

86     def parse_print_statement(self):
87         self.match("KEYWORD")
88         self.match("LPAREN")
89         arguments = []
90
91         while self.current_token() and self.current_token()[0] != "RPAREN":
92             if self.current_token()[0] == "COMMA":
93                 self.match("COMMA")
94             else:
95                 if self.current_token()[0] == "STRING":
96                     arguments.append(ParseNode(self.match("STRING")))
97                 elif self.current_token()[0] == "ID":
98                     arguments.append(ParseNode(self.match("ID")))
99                 else:
100                     raise SyntaxError(f"Unexpected token in print statement: {self.current_token()}")
101
102         self.match("RPAREN")
103         return ParseNode("Print", arguments)
104
105     def parse_expression(self):
106         left = self.parse_term()
107         while self.current_token() and self.current_token()[0] in {"PLUS", "MINUS"}:
108             op = self.match(self.current_token()[0])
109             right = self.parse_term()
110             left = ParseNode("Expression", [left, ParseNode(op), right])
111
112         while self.current_token() and self.current_token()[0] in {"LARGERTHAN", "LESSTHAN", "EQUALS", "NOTEQUALS"}:
113             op = self.match(self.current_token()[0])
114             right = self.parse_term()
115             left = ParseNode("RelationalExpression", [left, ParseNode(op), right])
116
117         return left
118
119     def parse_term(self):
120         left = self.parse_factor()
121         while self.current_token() and self.current_token()[0] in {"MULT", "DIV"}:
122             op = self.match(self.current_token()[0])
123             right = self.parse_factor()
124             left = ParseNode("Term", [left, ParseNode(op), right])
125         return left
126
127     def parse_factor(self):
128         token = self.current_token()
129         if token[0] == "NUMBER":
130             return ParseNode(self.match("NUMBER"))
131         elif token[0] == "ID":
132             return ParseNode(self.match("ID"))
133         elif token[0] == "STRING":
134             return ParseNode(self.match("STRING"))
135         elif token[0] == "LPAREN":
136             self.match("LPAREN")
137             expr = self.parse_expression()
138             self.match("RPAREN")
139             return expr
140         elif token[0] == "LBRACKET":
141             self.match("LBRACKET")
142             elements = []
143             while self.current_token() and self.current_token()[0] != "RBRACKET":
144                 if self.current_token()[0] == "COMMA":
145                     self.match("COMMA")
146                 else:
147                     elements.append(self.parse_expression())
148             self.match("RBRACKET")
149             return ParseNode("Array", elements)
150         else:
151             raise SyntaxError(f"Unexpected factor: {token}")
152
153

```

Produces a structured parse tree for the input tokens, suitable for further processing like interpretation or code generation.



How It Works

parse_program:

Starts parsing the entire program as a series of statements.

parse_statements:

Parses multiple statements (e.g., assignments, if blocks, print).

Specific Parsers:

parse_statement: Identifies and dispatches to the correct statement parser.

`parse_assignment`: Parses variable assignments (e.g., `x = 5`).

`parse_if_statement`: Parses if conditions, including optional else.

`parse_print_statement`: Parses print with arguments.

`parse_expression`: Handles arithmetic (+, -) and relational (<, ==) expressions.

`parse_term` / `parse_factor`: Parses more granular expressions, numbers, identifiers, and strings.

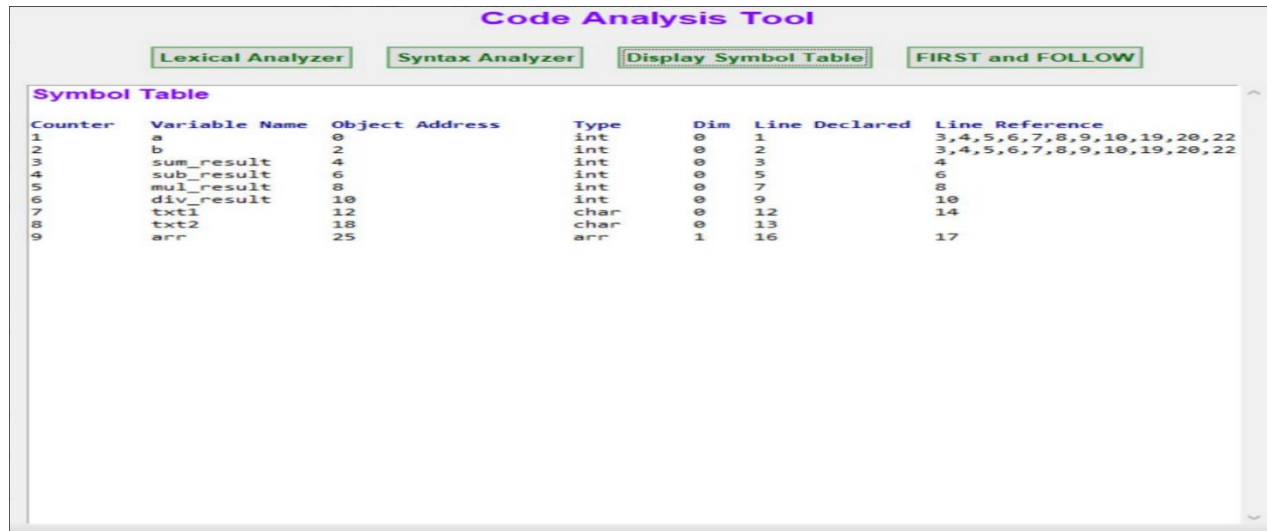
Error Handling:

Detects invalid tokens or mismatched structures (e.g., missing `:` or `)`).

symbol table generator : commonly used in compilers to store information about variables declared in the source code. It analyzes the code to extract metadata about variables, including their names, types, memory addresses, declaration, and references.

```
1 import re
2
3 TYPE_SIZES = {
4     'int': 2,
5     'float': 4,
6     'char': 1,
7 }
8
9 def generate_symbol_table(code):
10     pattern = r"(\w+)\s*=\s*(.*)"
11     lines = code.strip().split('\n')
12     symbol_table = []
13     current_address = 0
14
15     def get_var_type_and_size(var_value):
16         if re.match(r"^\d+\s*/\s*\(\s*\)\s*$", var_value):
17             return 'int', TYPE_SIZES['int']
18         elif var_value.startswith '[' and var_value.endswith(']'):
19             num_elements = len(re.findall(r'\d+', var_value))
20             return 'arr', TYPE_SIZES['int'] * num_elements
21         elif "'" in var_value:
22             return 'char', TYPE_SIZES['char'] * len(var_value.replace("'", ''))
23         else:
24             return 'int', TYPE_SIZES['int']
25
26     for line_num, line in enumerate(lines, 1):
27         match = re.match(pattern, line.strip())
28         if match:
29             var_name, var_value = match.groups()
30             var_type, var_size = get_var_type_and_size(var_value)
31             symbol_table.append({
32                 'counter': len(symbol_table) + 1,
33                 'Variable Name': var_name,
34                 'Object Address': current_address,
35                 'Type': var_type,
36                 'Dim': 0 if var_type in ['int', 'float', 'char'] else 1,
37                 'Line Declared': line_num,
38                 'Line Reference': []
39             })
40             current_address += var_size
41
42     for line_num, line in enumerate(lines, 1):
43         for var in symbol_table:
44             if var['Variable Name'] in line and var['Line Declared'] != line_num:
45                 if "print" in line.strip():
46                     if var['Variable Name'] in re.findall(r'\w+', line):
47                         var['Line Reference'].append(line_num)
48                 elif "[" not in line.strip():
49                     var['Line Reference'].append(line_num)
50
51     return symbol_table
52
```

Returns a list of dictionaries, where each dictionary represents a variable and its metadata.



The screenshot shows a window titled "Code Analysis Tool" with four buttons: "Lexical Analyzer", "Syntax Analyzer", "Display Symbol Table", and "FIRST and FOLLOW". The "Display Symbol Table" button is highlighted. Below the buttons is a table titled "Symbol Table" with the following data:

Counter	Variable Name	Object Address	Type	Dim	Line Declared	Line Reference
1	a	0	int	0	1	3,4,5,6,7,8,9,10,19,20,22
2	b	2	int	0	2	3,4,5,6,7,8,9,10,19,20,22
3	sum_result	4	int	0	3	4
4	sub_result	6	int	0	5	6
5	mul_result	8	int	0	7	8
6	div_result	10	int	0	9	10
7	txt1	12	char	0	12	14
8	txt2	18	char	0	13	14
9	arr	25	arr	1	16	17

? Input and Symbol Table Initialization:

Takes code (a multiline string with variable declarations).

Initializes symbol_table to store metadata about variables.

Starts current_address at 0 to calculate memory addresses.

? Type Detection (get_var_type_and_size):

Determines the type and size of variables based on their values:

Integer (int): Numeric expressions like `1 + 2`.

Array (arr): Square-bracketed values like `[1, 2, 3]`.

Character (char): Strings enclosed in quotes like `"hello"`.

? Variable Parsing:

Uses a regex pattern (`r"(\w+)\s*=\s*(.*)"`) to match variable assignments.

Extracts the variable name, value, type, size, and other metadata:

Dim: Dimensions (1 for arrays, 0 for scalars).

Line Declared: Line number where the variable is declared.

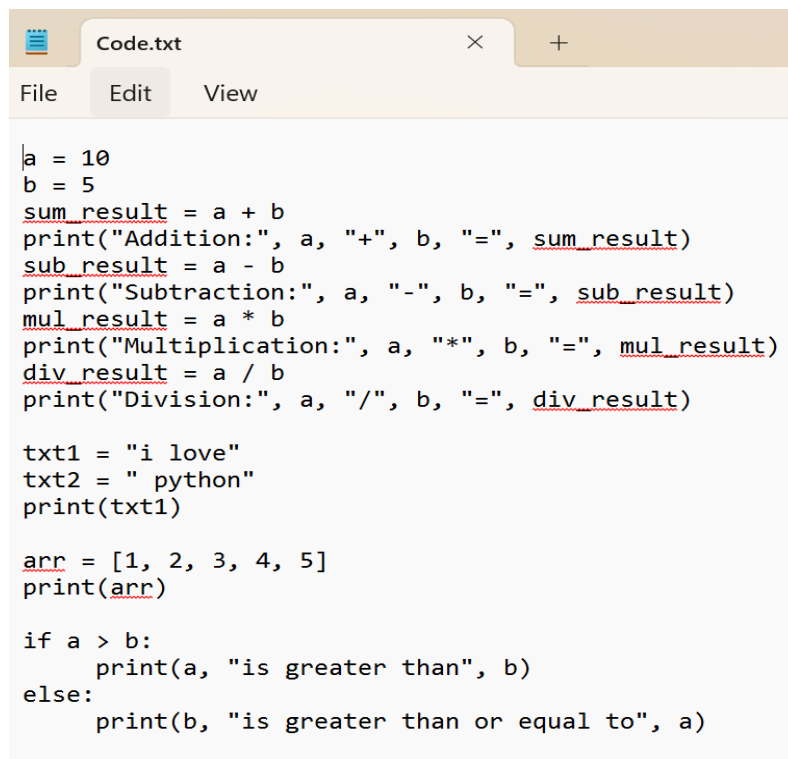
Updates `current_address` for each variable based on its size.

Reference Tracking:

For each line, checks if declared variables are referenced later in the code.

Updates Line Reference with line numbers where the variable is used.

the text file we used :



```
a = 10
b = 5
sum_result = a + b
print("Addition:", a, "+", b, "=", sum_result)
sub_result = a - b
print("Subtraction:", a, "-", b, "=", sub_result)
mul_result = a * b
print("Multiplication:", a, "*", b, "=", mul_result)
div_result = a / b
print("Division:", a, "/", b, "=", div_result)

txt1 = "i love"
txt2 = " python"
print(txt1)

arr = [1, 2, 3, 4, 5]
print(arr)

if a > b:
    print(a, "is greater than", b)
else:
    print(b, "is greater than or equal to", a)
```

first and follow :This code processes context-free grammar (CFG) rules to compute **FIRST** and **FOLLOW** sets while eliminating left recursion and performing left factoring. Here's a brief breakdown of its functionality:

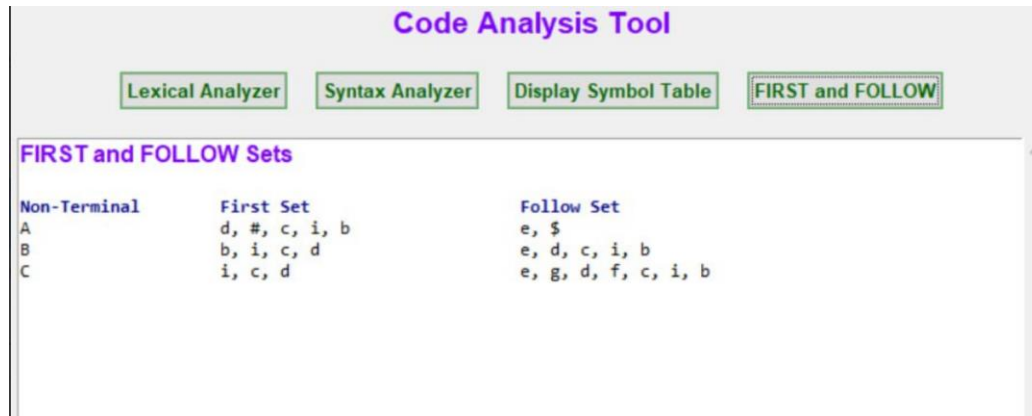
```
1 import re
2
3 start_symbol = ""
4 productions = {}
5 first_table = {}
6 follow_table = {}
7
8 def creatProduction(file_path):
9     global start_symbol, productions
10    with open(file_path, "r") as file:
11        for production in file:
12            lhs, rhs = re.split(r"->", production)
13            rhs = re.split(r"\\|\\n", rhs)
14            productions[lhs.strip()] = set(rhs) - {' '}
15            if not start_symbol:
16                start_symbol = lhs.strip()
17
18 def isNonterminal(symbol):
19     return symbol.isupper()
20
21 def eliminateLeftRecursion():
22     global productions
23     new_productions = {}
24     for nt, rules in productions.items():
25         alpha_rules = set()
26         beta_rules = set()
27         for rule in rules:
28             if rule.startswith(nt):
29                 alpha_rules.add(rule[len(nt):])
30             else:
31                 beta_rules.add(rule)
32
33         if alpha_rules:
34             nt_new = nt + ""
35             while nt_new in productions:
36                 nt_new += ""
37             new_productions[nt] = {beta + nt_new for beta in beta_rules}
38             new_productions[nt_new] = {alpha + nt_new for alpha in alpha_rules} | {"#"}
39         else:
40             new_productions[nt] = rules
41
42     productions = new_productions
43
44 def performLeftFactoring():
45     global productions
46     new_productions = {}
47     for nt, rules in productions.items():
48         prefix_map = {}
49         for rule in rules:
50             prefix = rule[0]
51             if prefix not in prefix_map:
52                 prefix_map[prefix] = []
53             prefix_map[prefix].append(rule)
54
55         new_rules = set()
56         for prefix, grouped_rules in prefix_map.items():
57             if len(grouped_rules) > 1:
58                 nt_new = nt + ""
59                 while nt_new in productions:
60                     nt_new += ""
61                 new_rules.add(prefix + nt_new)
62                 new_productions[nt_new] = {rule[1:] or "#" for rule in grouped_rules}
63             else:
64                 new_rules.add(grouped_rules[0])
65         new_productions[nt] = new_rules
66
67     productions.update(new_productions)
68
69 def firstFunc(symbol):
70     if symbol in first_table:
71         return first_table[symbol]
72
73     if isNonterminal(symbol):
74         first = set()
75         for production in productions[symbol]:
76             if production:
77                 first_element = production[0]
78                 fst = firstFunc(first_element)
79                 first = first.union(fst)
80     return first
81
82 else:
83     return set(symbol)
```

```

84 def followFunc(symbol):
85     if symbol not in follow_table:
86         follow_table[symbol] = set()
87     for nt in productions.keys():
88         for rule in productions[nt]:
89             pos = rule.find(symbol)
90             if pos != -1:
91                 if pos == (len(rule) - 1):
92                     if nt != symbol:
93                         follow_table[symbol] = follow_table[symbol].union(followFunc(nt))
94                 else:
95                     first_next = set()
96                     for next in rule[pos + 1:]:
97                         fst_next = firstFunc(next)
98                         first_next = first_next.union(fst_next - {'#'})
99                         if '#' not in fst_next:
100                             break
101                     if '#' in fst_next:
102                         if nt != symbol:
103                             follow_table[symbol] = follow_table[symbol].union(followFunc(nt))
104                             follow_table[symbol] = follow_table[symbol].union(first_next) - {'#'}
105                     else:
106                         follow_table[symbol] = follow_table[symbol].union(first_next)
107     return follow_table[symbol]
108
109 def compute_first_follow(file_path):
110     global first_table, follow_table
111     creatProduction(file_path)
112     eliminateLeftRecursion()
113     performLeftFactoring()
114
115     for nt in productions:
116         first_table[nt] = firstFunc(nt)
117     follow_table[start_symbol] = set('$')
118     for nt in productions:
119         follow_table[nt] = followFunc(nt)
120
121     result = {}
122     for nt in productions:
123         result[nt] = (list(first_table[nt]), list(follow_table[nt]))
124     return result
125
126

```

Produces a dictionary with non-terminals as keys, and their respective **FIRST** and **FOLLOW** sets as values.



The screenshot shows a web application titled "Code Analysis Tool" with four buttons: "Lexical Analyzer", "Syntax Analyzer", "Display Symbol Table", and "FIRST and FOLLOW". The "FIRST and FOLLOW" button is highlighted. Below the buttons, the section "FIRST and FOLLOW Sets" displays a table with three columns: "Non-Terminal", "First Set", and "Follow Set".

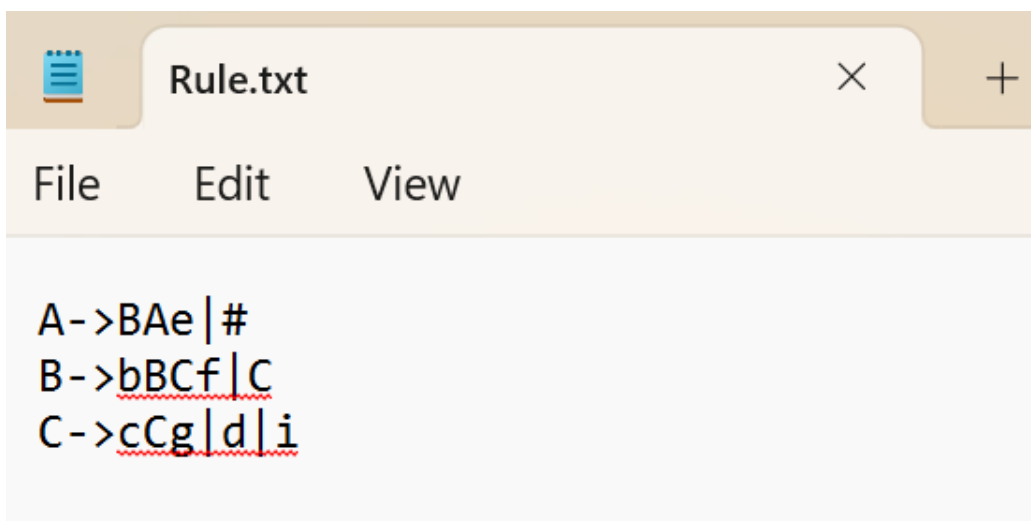
Non-Terminal	First Set	Follow Set
A	d, #, c, i, b	e, \$
B	b, i, c, d	e, d, c, i, b
C	i, c, d	e, g, d, f, c, i, b

Key Features:

1. **Input:** Reads CFG from a file where each line defines a production, e.g., S
-> Aa | b.
2. **Elimination of Left Recursion:**
 - Identifies rules with left recursion (e.g., A -> Aa | b).
 - Rewrites them into non-left-recursive form using auxiliary non-terminals.
3. **Left Factoring:**
 - Detects common prefixes in production rules (e.g., A -> ab | ac).
 - Refactors them to remove ambiguity (e.g., A -> aA', A' -> b | c).
4. **FIRST Function:**
 - Computes the set of terminal symbols that can appear at the start of strings derived from a given symbol.
5. **FOLLOW Function:**

- Computes the set of terminal symbols that can appear immediately after a given non-terminal in any derivation.
- Starts with $\text{FOLLOW}(\text{start_symbol}) = \{\text{'\$'}\}$ (end-of-input marker).

the rule we used is :



```
Rule.txt
File Edit View
A->BAe|#
B->bBCf|C
C->cCg|d|i
```

The Mian code :

```
8 def analyze_lexical():
9     with open("Code.txt", "r") as file:
10         code = file.read()
11
12     result_text.delete(1.0, tk.END)
13     result_text.insert(tk.END, "Lexical Analysis (Scanner)\n\n", "header")
14     tokens = tokenize(code)
15     result_text.insert(tk.END, "Tokens:\n", "subheader")
16     result_text.insert(tk.END, "\n\n")
17     result_text.insert(tk.END, str(tokens) + "\n\n")
18
19     result_text.tag_configure("header", foreground="87A09F", font=("Arial", 12, "bold"))
20     result_text.tag_configure("subheader", foreground="darkgreen", font=("Arial", 11, "italic"))
21
22 def analyze_parser():
23     with open("Code.txt", "r") as file:
24         code = file.read()
25
26     result_text.delete(1.0, tk.END)
27     result_text.insert(tk.END, "Syntax Analysis (Parser)\n\n", "header")
28     tokens = tokenize(code)
29     parser = Parser(tokens)
30     parse_tree = parser.parse_program()
31     result_text.insert(tk.END, "Parse Tree:\n", "subheader")
32     result_text.insert(tk.END, "\n\n")
33     result_text.insert(tk.END, str(parse_tree) + "\n\n", "indented")
34
35     result_text.tag_configure("header", foreground="87A09F", font=("Arial", 12, "bold"))
36     result_text.tag_configure("subheader", foreground="darkgreen", font=("Arial", 11, "italic"))
37     result_text.tag_configure("indented", indent=20, font=("Consolas", 10))
38
39 def display_symbol_table():
40     with open("Code.txt", "r") as file:
41         code = file.read()
42
43     symbol_table = generate_symbol_table(code)
44
45     result_text.delete(1.0, tk.END)
46     result_text.insert(tk.END, "Symbol Table\n\n", "header")
47     result_text.insert(tk.END, f"{'Counter':<10}{'Variable Name':<15}{'Object Address':<20}{'Type':<10}{'Din':<5}{'Line Declared':<15}{'Line Reference':<15}\n", "table_header")
48
49     for entry in symbol_table:
50         line = f"{entry['counter']:<10}{entry['Variable Name']:<15}{entry['Object Address']:<20}{entry['Type']:<10}{entry['Din']:<5}{entry['Line Declared']:<15}{', '.join(map(str, entry['Line Reference']))}\n"
51         result_text.insert(tk.END, line)
52
53     result_text.tag_configure("header", foreground="87A09F", font=("Arial", 12, "bold"))
54     result_text.tag_configure("table_header", foreground="darkblue", font=("Consolas", 10, "bold"))
55
56 def display_first_follow():
57     result_text.delete(1.0, tk.END)
58     result_text.insert(tk.END, "FIRST and FOLLOW Sets\n\n", "header")
59
60     first_follow = compute_first_follow("Rule.txt")
61
62     result_text.insert(tk.END, f"{'Non-Terminal':<20}{'First Set':<30}{'Follow Set':<30}\n", "table_header")
63     for nt, (first_set, follow_set) in first_follow.items():
64         line = f"{nt:<20}{', '.join(first_set)<30}{', '.join(follow_set)}\n"
65         result_text.insert(tk.END, line)
66
67     result_text.tag_configure("header", foreground="87A09F", font=("Arial", 12, "bold"))
68     result_text.tag_configure("table_header", foreground="darkblue", font=("Consolas", 10, "bold"))
69
70 root = tk.Tk()
71 root.title("Compiler Design")
72 root.geometry("900x700")
73
74 frame = ttk.Frame(root, padding=(40, 0, 0, 0))
75 frame.grid(row=0, column=0, sticky=(tk.W, tk.E, tk.S, tk.N), tk.S)
76
77 header_label = ttk.Label(frame, text="Code Analysis Tool", font=("Arial", 16, "bold"), foreground="87A09F")
78 header_label.grid(row=0, column=1, pady=10)
79
80 style = ttk.Style()
81 style.configure(
82     "Custom.TButton",
83     background="black",
84     foreground="darkgreen",
85     height=35,
86     font=("Arial", 10, "bold")
87 )
88
89 button_frame = ttk.Frame(frame)
90 button_frame.grid(row=1, column=0, columnspan=3, pady=10, padx=10)
91
92 lexical_button = ttk.Button(button_frame, text="Lexical Analyzer", command=analyze_lexical, style="Custom.TButton")
93 lexical_button.grid(row=0, column=0, padx=10)
94
95 parser_button = ttk.Button(button_frame, text="Syntax Analyzer", command=analyze_parser, style="Custom.TButton")
96 parser_button.grid(row=0, column=1, padx=10)
97
98 symbol_table_button = ttk.Button(button_frame, text="Display Symbol Table", command=display_symbol_table, style="Custom.TButton")
99 symbol_table_button.grid(row=0, column=2, padx=10)
100
101 first_follow_button = ttk.Button(button_frame, text="FIRST and FOLLOW", command=display_first_follow, style="Custom.TButton")
102 first_follow_button.grid(row=0, column=3, padx=10)
103
104 result_text = scrolledtext.ScrolledText(frame, wrap=tk.WORD, width=100, height=15, font=("Consolas", 10))
105 result_text.grid(row=2, column=0, columnspan=3, pady=10)
106
107 root.mainloop()
108
```

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