Designing Schema for Hive



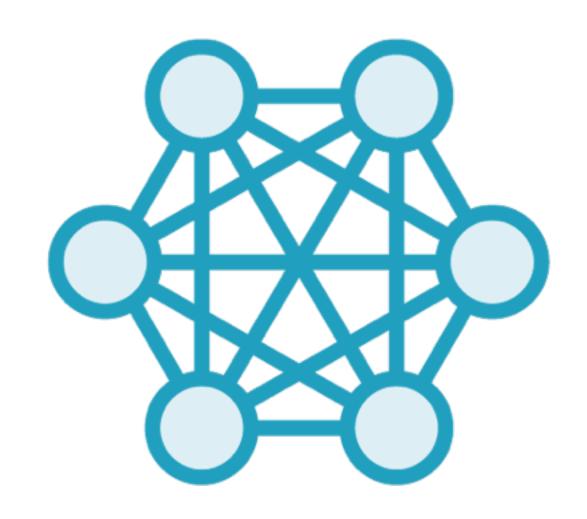
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Overview

Understand normalized storage in traditional databases

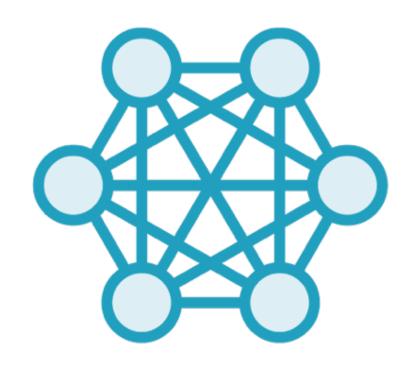
Understand why Hive stores data in denormalized form

Hive Schema Design



Complex process

Based on the type of dataset and kind of queries that are run



Column structure and data types

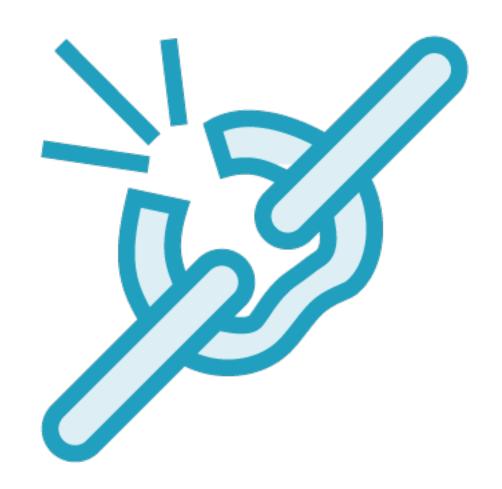
Serialized forms in HDFS

Partitioning to improve query performance

Bucketing to tweak sampling and joins

Normalized Storage in Traditional Databases

Traditional Database Design



Normalized data

Data is stored in a granular form to minimize redundancy

Employee Information

name

address

id

subordinates

department

grade

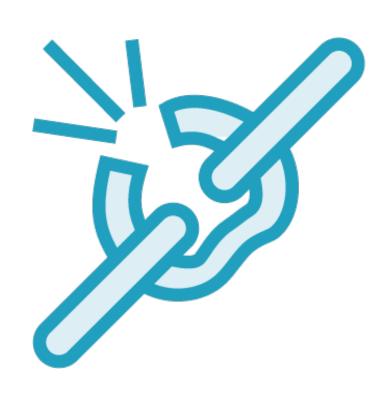
id name grade department



id subordinates

id address

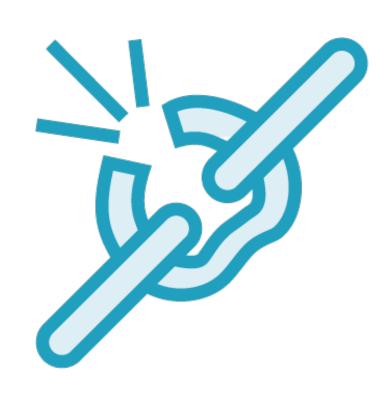
Minimize Redundancy



Employee Details

Employee Subordinates

Employee Address



Employee Details

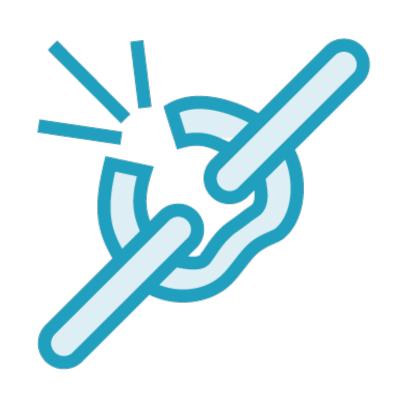
ld	Name	Department	Grade
1	Emily	Finance	6

Employee Subordinates

ld	Subordinate Id
1	2
1	3

Employee Address

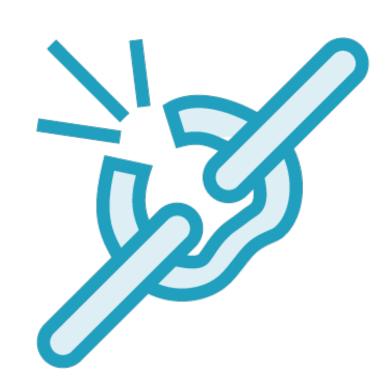
ld	City	Zip Code
1	Palo Alto	94305
2	Seattle	98101



Employee Details

ld	Name	Function	Grade
1	Emily	Finance	6
2	John	Finance	3
3	Ben	Finance	4

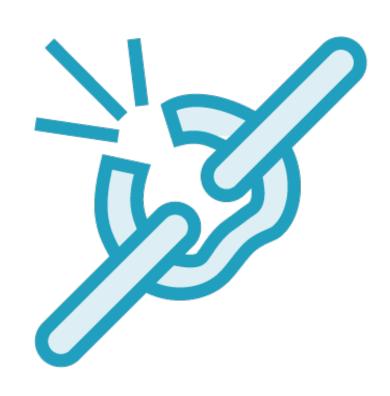
All employee details in one table



Employee Subordinates

ia	Subordinate id
1	2
1	3

Employees referenced only by ids everywhere else



Employee Address

ld	City	Zip Code
1	Palo Alto	94305
2	Seattle	98101

Data is made more granular by splitting it across multiple tables

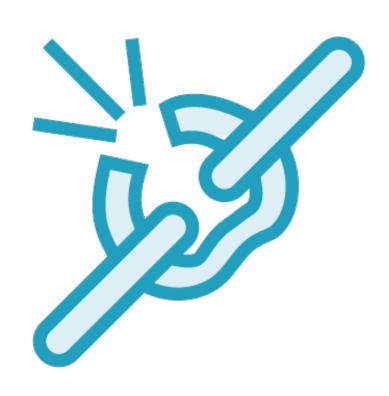


ld	Name	Function	Grade
1	Emily	Finance	6

ld	Subordinate Id
1	2
1	3

ld	City	Zip Code
1	Palo Alto	94305
2	Seattle	98101

Normalization



ld	Name	Function	Grade
1	Emily	Finance	6

join

ld	Subordinate Id
1	2
1	3

Query for Emily's department and her subordinates

Normalization

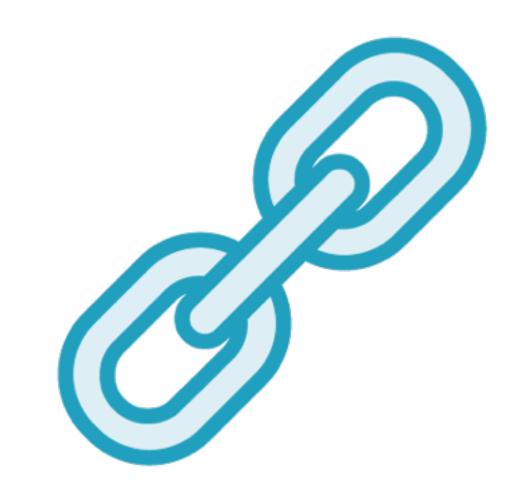


Minimizes redundancy, optimizes storage

Foreign keys to ensure valid joins

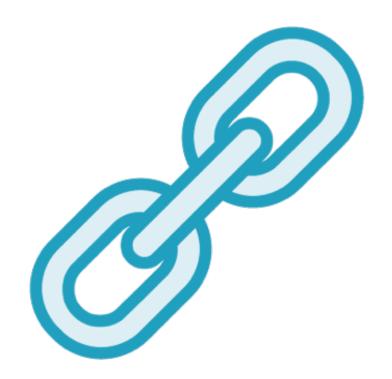
Updates in one location, no duplication of data

Denormalized Storage in Hive



Denormalized data

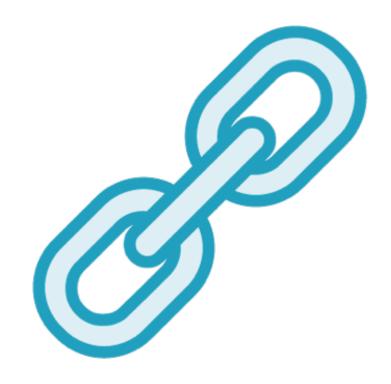
Data is compressed into one table to be read in a single operation



Disk space is very cheap

No foreign key constraints

Read operations, no data updates



Optimize the number of disk seeks

Store data for an entity in one location

Ignore redundancy, minimize joins

ld	Name	Function	Grade
1	Emily	Finance	6
2	John	Finance	3
3	Ben	Finance	4

ld	Subordinate Id
1	2
1	3



ld	Name	Function	Grade
1	Emily	Finance	6
2	John	Finance	3
3	Ben	Finance	4

Id	Subordinate Id	
1	2	
1	3	



ld	Name	Function	Grade	Subordinates
1	Emily	Finance	6	<array></array>
2	John	Finance	3	
3	Ben	Finance	4	

ld	Name	Function	Grade
1	Emily	Finance	6
2	John	Finance	3
3	Ben	Finance	4

ld	Subordinate Id	
1	2	
1	3	
	1	



ld	Name	Function	Grade	Subordinates
1	Emily	Finance	6	2,3
2	John	Finance	3	
3	Ben	Finance	4	

ld	Name	Function	Grade
1	Emily	Finance	6
2	John	Finance	3
3	Ben	Finance	4

ld	City	Zip Code	
1	Palo Alto	94305	
2	Seattle	98101	



ld	Name	Function	Grade	Subordinates	Address
1	Emily	Finance	6	2,3	<struct></struct>
2	John	Finance	3		
3	Ben	Finance	4		

ld	Name	Function	Grade
1	Emily	Finance	6
2	John	Finance	3
3	Ben	Finance	4

ld	City	Zip Code	
1	Palo Alto	94305	
2	Seattle	98101	



ld	Name	Function	Grade	Subordinates	Address
1	Emily	Finance	6	2,3	Palo Alto, 94305
2	John	Finance	3		
3	Ben	Finance	4		

ld	Name	Function	Grade	Subordinates	Address
1	Emily	Finance	6	2,3	Palo Alto, 94305
2	John	Finance	3		
3	Ben	Finance	4		

Store everything related to an employee in the same table

ld	Name	Function	Grade	Subordinates	Address
1	Emily	Finance	6	2,3	Palo Alto, 94305
2	John	Finance	3		
3	Ben	Finance	4		

Get all details about an employee in one read operation

Summary

Understood the differences between normalized and denormalized storage

Understood why Hive stores data in denormalized form