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10 ARM Machine Code: Branching and Memory Instruc- tions

CPE 221

The University of Alabama in Huntsville

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Table of Content

Memory Instruction

Branching Instruction

Encoding PUSH and POP Instruction

Memory Instruction

1	0	1	1	1	1	0	1	1	0	0	1	0	0	0	0	0	1	0	1
1	1	1	0	1	1	0	0	1	0	0	1	0	1	0	1	0	0	0	0
1	1	0	0	1	0	0	0	1	0	1	1	1	0	0	0	0	0	1	1

Memory Instructions

31:28	27:26	25	24	23	22	21	20	19:16	15:12	11:0
cond	op	\bar{I}	P	U	B	W	L	Rn	Rd	Src2
4 bits	2 bits	1 bit	1 bit	1 bit	1 bit	1 bit	1 bit	4 bits	4 bits	12 bits

If $\bar{I} = 1$ (no immediate) but with register shifting:

Bits 11:7	Bits 6:5	Bit 4	Bits 3:0
shamt5	sh	1/0	Rm
5-bit shift amount	specific shift instructions	1 in the text-book, 0 in the copulator	Register used for shifting

Memory Instructions

31:28	27:26	25	24	23	22	21	20	19:16	15:12	11:0
cond	op	\bar{I}	P	U	B	W	L	Rn	Rd	Src2
4 bits	2 bits	1 bit	1 bit	1 bit	1 bit	1 bit	1 bit	4 bits	4 bits	12 bits

If $\bar{I} = 1$ (no immediate) no shifting, just a register:

Bits 11:7	Bits 6:5	Bit 4	Bits 3:0
shamt5	sh	0	Rm
00000	00	0 (default)	Register

Memory Instructions

31:28	27:26	25	24	23	22	21	20	19:16	15:12	11:0
cond	op	\bar{I}	P	U	B	W	L	Rn	Rd	Src2
4 bits	2 bits	1 bit	1 bit	1 bit	1 bit	1 bit	1 bit	4 bits	4 bits	12 bits

If $\bar{I} = 0$ (immediate):

Bits 11:0
imm12
12 bit immediate

Memory Instructions

31:28	27:26	25	24	23	22	21	20	19:16	15:12	11:0
cond	op	\bar{I}	P	U	B	W	L	Rn	Rd	Src2
4 bits	2 bits	1 bit	1 bit	1 bit	1 bit	1 bit	1 bit	4 bits	4 bits	12 bits

I = Immediate, P = Pre-index, U = Add, W = Writeback, L = Load, B = Byte

Memory Instructions

31:28	27:26	25	24	23	22	21	20	19:16	15:12	11:0
cond	op	\bar{I}	P	U	B	W	L	Rn	Rd	Src2
4 bits	2 bits	1 bit	1 bit	1 bit	1 bit	1 bit	1 bit	4 bits	4 bits	12 bits

P	W	Index-Mode
0	0	Post-index
0	1	Not Supported
1	0	Offset
1	1	Pre-index

Memory Instructions

31:28	27:26	25	24	23	22	21	20	19:16	15:12	11:0
cond	op	\bar{I}	P	U	B	W	L	Rn	Rd	Src2
4 bits	2 bits	1 bit	1 bit	1 bit	1 bit	1 bit	1 bit	4 bits	4 bits	12 bits

L	B	Instruction
0	0	STR
0	1	STRB
1	0	LDR
1	1	LDRB

Example 1

Encode STR R11, [R5], #-26

This instruction has indexing type: post-indexing: $P = 0$, $W = 0$.
Immediate offset is subtracted from the base. So $I = 0$, $U = 0$.

31:28	27:26	25	24	23	22	21	20	19:16	15:12	11:0
cond	op	\bar{I}	P	U	B	W	L	Rn	Rd	Src2
1110	01	0	0	0	0	0	0	0101	1011	0000 0001 1010

0x E405B017

Example 2

Encode LDR R2, [R3, #16] !

Pre-indexed Word Load:

- ▶ Load word into R2, $L = 1$, $B = 0$ (no Byte)
- ▶ Base register R3
- ▶ Immediate offset of 16, which is positive, $U = 1$
- ▶ Pre-indexed (update base register) $P = 1$ $W = 1$, P

31:28	27:26	25	24	23	22	21	20	19:16	15:12	11:0
cond	op	\bar{I}	P	U	B	W	L	Rn	Rd	Src2
1110	01	0	1	1	0	1	1	0011	0010	0000 0001 0000

0x E5B32010

Example 3

Encode STRB R7, [R4], R6

Post-indexed Byte Store:

- ▶ Store byte from R7, L = 0, Byte, hence B = 1
- ▶ Base register R4
- ▶ Offset from R6 register, no Immediate, sets U = 1
- ▶ Post-indexed, P = 0, W = 0

31:28	27:26	25	24	23	22	21	20	19:16	15:12	11:0
cond	op	\bar{I}	P	U	B	W	L	Rn	Rd	Src2
1110	01	1	0	1	1	0	0	0100	0111	0000 0000 0110

0x E6C47006

Example 4

Encode LDR R11, [R7, R8, LSL #2]

Word Load with Register Shift Offset:

- ▶ Load word into R11, $L = 1$, $B = 0$
- ▶ Base register R7
- ▶ Offset from R8, logically shifted left by 2, $U = 1$, Offset, so $P = 1$, $W = 0$
- ▶ Register-based offset with shift, $sh = 00$ for LSL

31:28	27:26	25	24	23	22	21	20	19:16	15:12	11:0
cond	op	\bar{I}	P	U	B	W	L	Rn	Rd	Src2
1110	01	1	1	1	0	0	1	0111	1011	00010 00 0 1000

0x E797B108

Branching Instruction

0	0	1	1	1	1	0	1	1	1	0	0	1	0	0	0	0	0	1	0
1	1	0	1	1	0	0	1	0	0	1	1	0	0	0	1	0	0	1	0
1	1	1	0	1	0	1	1	0	1	0	1	0	0	1	1	0	0	1	1

Encoding Branching Instructions: B Label

31:28	27:26	25	24	23:0
cond	op	1	L	imm24
cond	10	1	$L = 0$	imm24 is in 2's complement, used to specify an instruction address relative to $PC + 8$.

Encoding Branching Instructions: BL Label

Branch with Link

31:28	27:26	25	24	23:0
cond	op	1	L	imm24
cond	10	1	$L = 1$	imm24 is in 2's complement, used to specify an instruction address relative to $PC + 8$.

Example 5

Encode

0x80A0 BLT THERE

When we have the rest of the program as:

0x80A4 ADD R0, R1, R2

0x80A8 SUB R0, R0, R9

0x80AC ADD SP, SP, #8

0x80B0 MOV PC, LR

0x80B4 THERE SUB R0, R0, #1

0x80B8 ADD R3, R3, #0x5

Example 5

Step 1: We calculate BTA. BTA = Branch Target Address, the instruction address to be executed when the branch is taken. BLT has BTA of 0x80B4, the instruction address of THERE.

Step 2: 24-bit immediate field gives the number of instructions between the BTA and PC+8 (two instructions past the branch)

imm24 here is 3, as BTA is three instructions past PC + 8 (0x80A8).

The processor calculates BTA from the instruction by sign-extending the 24-bit immediate, shifting it left 2 (to convert word to bytes) and adding it to PC-8.

31:28	27:26	25	24	23:0
cond	op	1	L	imm24
1011	10	1	0	0000 0000 0000 0000 0000 0011

Example 6

If the instruction at address 0x8400 is 0xEA000005, what is the branch target address?

Solution:

- ▶ This is an unconditional branch (B) instruction
- ▶ The immediate field is 5
- ▶ Branch target = $PC+8 + (\text{imm}_{24} \times 4)$
- ▶ $PC+8$ for the instruction at 0x8400 is 0x8408
- ▶ $5 \times 4 = 20$ (0x14)
- ▶ $0x8408 + 0x14 = 0x841C$

Conditions Encoding

cond	Mnemonic	Name	CondEx
0000	EQ	Equal	Z
0001	NE	Not equal	\bar{Z}
0010	CS / HS	Carry set / unsigned higher or same	C
0011	CC / LO	Carry clear / unsigned lower	\bar{C}
0100	MI	Minus / negative	N
0101	PL	Plus / positive or zero	\bar{N}
0110	VS	Overflow / overflow set	V
0111	VC	No overflow / overflow clear	\bar{V}
1000	HI	Unsigned higher	$\bar{Z}C$
1001	LS	Unsigned lower or same	$Z \text{ OR } \bar{C}$
1010	GE	Signed greater than or equal	$\bar{N} \oplus \bar{V}$
1011	LT	Signed less than	$N \oplus V$
1100	GT	Signed greater than	$\bar{Z} (\bar{N} \oplus V)$
1101	LE	Signed less than or equal	$Z \text{ OR } (N \oplus V)$
1110	AL (or none)	Always / unconditional	Ignored

Encoding PUSH and POP Instruction

1	1	0	0	1	1	0	0	0	0	0	1	1	1	1	0	0	1	0	0
1	1	0	1	0	0	0	1	0	1	1	1	1	1	0	0	0	1	0	1
1	0	0	0	0	1	0	0	1	0	0	0	1	0	0	0	1	0	0	1

PUSH Instruction

PUSH {R4, R5, LR}

is equivalent to

STMFD SP! {R4, R5, LR}

31:28	27:26	25	24	23	22	21	20	19:16	15:0
cond	op	I	P	U	S	W	L	Rn	register list
1110	10	0	1	0	0	1	0	1101	0100000000110000

P = 1 Pre-indexing mode, U = 0 for decrement, W = 1 write back, updates SP after storing, L = 0 for STM (Store operation)

Rn is the base register which is SP. Register list uses bit masking. Hence, encoding is 0x E92D4030.

POP Instruction

POP {R4, R5, LR}

is equivalent to

LDMFD SP! {R4, R5, LR}

31:28	27:26	25	24	23	22	21	20	19:16	15:0
cond	op	I	P	U	S	W	L	Rn	register list
1110	10	0	0	1	0	1	1	1101	0100000000110000

P = 0 Post-indexing mode, U = 1 for increment, W = 1 write back, updates SP after loading, L = 1 for LDM (Load operation)

Rn is the base register which is SP. Register list uses bit masking. Hence, encoding is 0x E8BD4030.

The End