

$$\begin{array}{r} a = 0 \\ b = 1 \end{array} \Rightarrow \begin{array}{r} 0000 \\ 0001 \\ \hline 0001 \end{array} \Rightarrow \textcircled{1}$$

$$\begin{array}{r} a = 1 \Rightarrow 0001 \\ b = 1 \Rightarrow 0001 \\ \hline 0010 \end{array} \Rightarrow \textcircled{2}$$

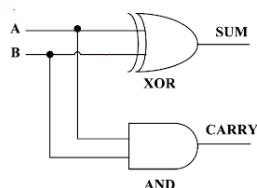
$$\begin{array}{r} a=3 \\ d=5 \end{array} \Rightarrow \begin{array}{r} 11 \\ 0011 \\ 0101 \\ \hline 1000 \end{array} \Rightarrow \textcircled{3} \times$$

$$\begin{array}{r} 0011 \\ 0101 \\ \hline 0110 \end{array} \quad \begin{array}{r} 0011 \\ 0101 \\ \hline 0001 \end{array}$$

sum \oplus carry \otimes

\oplus ADD + XOR

$$\begin{array}{r} 0110 \\ + 0010 \\ \hline 10 \end{array}$$



$$\begin{array}{r} 0010 \\ + 0010 \\ \hline 0100 \end{array}$$

Sum $\textcircled{1}$ \oplus carry \otimes

0100

$$\begin{array}{r} 0100 \\ \textcircled{1} \quad \textcircled{2} \\ \hline 0000 \end{array}$$

0000 1000

$$\begin{array}{r} 0000 \quad 0000 \\ 0000 \quad 1000 \\ \hline 0000 \quad 1000 \end{array} \Rightarrow \text{sum} \times$$

1 | 0000 1000 \oplus 0000 \otimes

$$\begin{array}{r} 0000 \quad 0000 \\ 0000 \quad 0000 \\ \hline 0000 \quad 0000 \end{array} \Rightarrow \textcircled{0} \times$$

(a,b)

... 111 111 111

while ($a \neq 0$) {

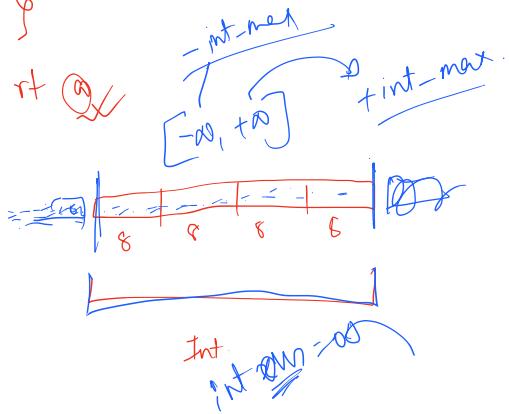
sum = a ^ b

carry = a & b

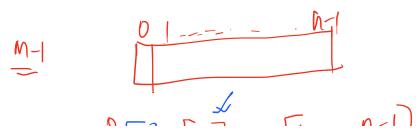
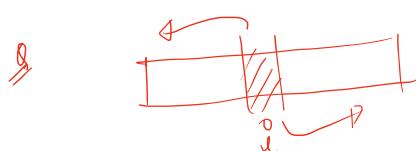
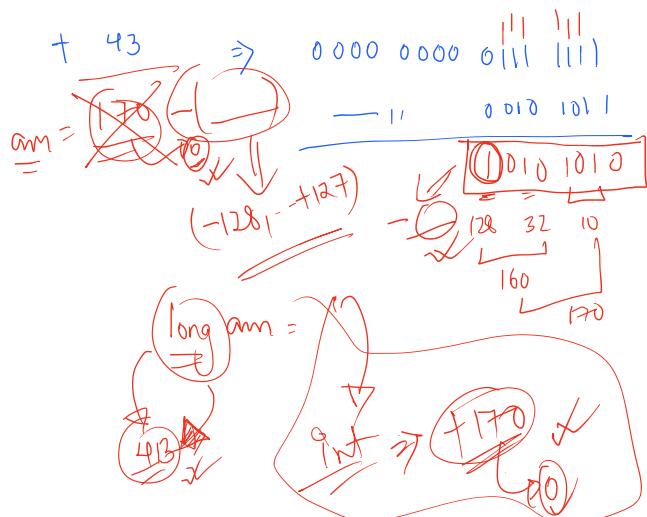
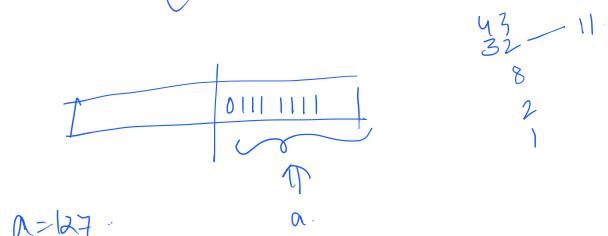
a = sum

b = (carry << 1)

rt



$$\text{byte} \Rightarrow 8 \text{ bits} \Rightarrow \begin{cases} -2^7 \Rightarrow -128 \\ +2^7 - 1 \Rightarrow +127 \end{cases}$$



for $i = 0 \dots n-1$ {

lsm = 0

$$\text{for } (j=0; j \leq i; j++)$$

$Lsm = arr[j]$

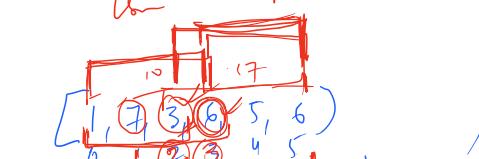
$$Lsm = arr[j]$$

$Rsm = arr[i+1]$

$$i=2 \quad [0, 1, 2] \quad [2 \dots n-1]$$

Lsm

Rsm



$i=0$	Lsm	Rsm
	0	(2)
	1	(3)
	2	(4)
	3	(5)
	4	(6)
	5	(7)

for $(j=0; j \leq i; j++)$

$Lsm = arr[j]$

$Rsm \rightarrow$

for $(j=i; j \leq n-1; j++)$

$Rsm = arr[j]$

if $(Lsm < Rsm)$

$\not\equiv 0$

$fret0$

i	0	1	2	3	4	5
	0	1	2	3	4	5

$\sum(0, i) \Rightarrow \underline{\underline{0(i)}}$

$$\sum(0, i); \quad \begin{array}{l} i \neq 0 \\ \text{and } 0 \leq i \end{array} \Rightarrow \sum(1, 3) \Rightarrow [1] + [2] + [3]$$

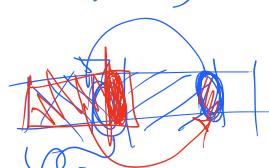
\downarrow

$\boxed{\text{fret}(j) - \text{fret}(i)}$

$\Rightarrow \text{fret}(3) - \text{fret}(0)$

* $\sum(2, 4) \geq [2] + [3] + [4]$

\downarrow



for $(i=0; i < n; i++)$

$Lsm = \text{fret}(i)$

// $\sum(i, n-1)$

\downarrow

$\dots 1 \ 0 \ \dots n$

$\left\{ \sum(0, i) \right\}$

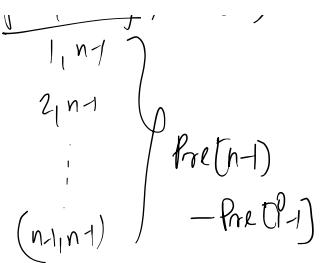
$0, 0, 1, 0, 2, \dots, (0, n-1)$

$\sum(0, n-1) \Rightarrow \text{fret}(n-1)$

$\sum(0, n-1) \Rightarrow \text{fret}(n-1)$

$$\text{if } (i > n) \\ R_{\text{sum}} = \text{Pre}[n-1]$$

$$\text{else} \\ R_{\text{sum}} = \text{Pre}[n-1] - \text{Pre}[i-1]$$



Pre	1	8	11	17	22	28
	0	1	2	3	4	5

i	L _{sum}	R _{sum}	
0	1	Pre[1:n-1] 28	X
1	8	28 - 1 \Rightarrow 27	X
2	11	28 - 8 = 20	X
3	17	28 - 11 \Rightarrow 17	✓

// prefix[0] = new $\text{int}(n)$

prefix[0] = arr[0]

for ($i=1; i < n; i++$) {

 prefix[i] = prefix[i-1] + arr[i]

}

$O(n^2)$

$\Rightarrow O(n)$

$$\text{TC} \Rightarrow \underbrace{O(n)}_{\text{Pre}} + \underbrace{O(n)}_{\text{Cal.}} \Rightarrow O(n^2)$$

$\therefore O(n) \approx$

1	7	3	6	5	6
0	1	2	3	4	5

$$R_{\text{sum}} = (28)$$

i	L _{sum}	R _{sum}	
0	1	28	X
1	8	28 - 1 \Rightarrow 27	X
2	11	28 - 8 - 1 \Rightarrow 21	X

(S)x ~ O(n)

$f_{\text{sum}} = 0$

for ($i=0; i < n; i++$) { }
 $f_{\text{sum}} += \text{array}[i]$

$L_{\text{sum}} = 0$
for ($i=0; i < n; i++$) { }
 $L_{\text{sum}} += \text{array}[i]$
if ($L_{\text{sum}} == f_{\text{sum}}$) E $O(n)$
if (i) ✗ X

$f_{\text{sum}} = \text{array}[0]$ TC = $O(1)$
at Θ ✓ SC = $O(1)$

$O(n^2) \xrightarrow{\text{Dra}} O(n)$ $O(n) \xrightarrow{\text{O}(n)} \Rightarrow \boxed{O(n)}$ ✗
 $O(1) \xrightarrow{\text{X}} O(1)$