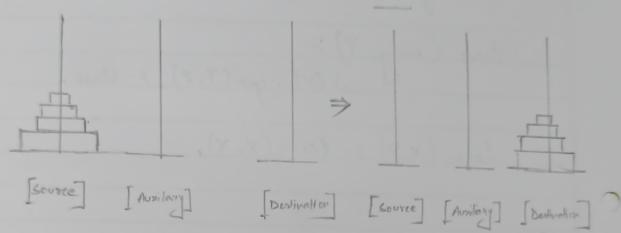
Expt No.	Page No Date
	g. Facts, Fully is a cat. Fully has block epots. Agaro is a log. Figure has white epots.
	Rules: Many owns a pet if it is a cat and it has
Œ	Suple-ment the rule base in CWP Prologo
	Algorithm /
	Flow Charet of John Charles of
	(in) Then Fubby has black spots (iv) Agover has white spots.
(Ç	(vi) Then check if it is a cat it has black operation of heart she laws that
	Programme listing !-
	Cat (fully), dog (figares),
Sreehari	Spot (fully, black),
	Teacher's Signature

	r
xpt No.	Page No Date
	spot (figure, white),
	Owne (mary, Υ):- cat(Υ), spot(Υ , Z), 2 = black.
E	Lones (X,Y) :- LOWIS (X,Y),
	Output: $Cat(X) \rightarrow x = fubby$
-	spot (ferbly), x) -> x -> black spot (framen, x) -> x -> white.
-	Owne (Mary, fubby) > True.
2 -	loves (Mary, Y) => Y = fully
-	
-	haden of a stay profession and a
EHARI _	
	Teacher's Signature

Tower of Hanoi

Source to Dartination Convention.



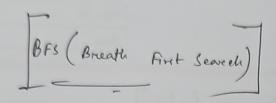
Output of Python Preogramme:

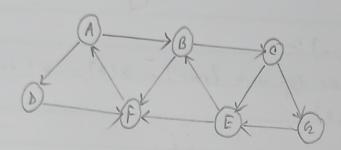
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É	In pleasant the Tower of Hanoi cling ent prealog & python?
	deparithme: - (i) Nous ut diske from source to auxiliary
_	Mone the NHA disk from Source to destination (11) Mone u-1 links from auxiliary to destination
6 -	START -> Precedence TOH (disk, sonwere, dert, and) If (disk == 1) Then, more disk from somee to destination
-	FISE, TOH (dick -1, sources, dest aux, dest)
-	TOH (dick-1, aux, dest, source)
-	End IF, END Procedure
-	Stop
•	Progream listing of Tower of Hanoi implementation in Pythou
	3. del ton (n, souver, dent, and) if (n==1): print (" Mone dick 1 from tower", sowere "to tower") retween
	print 1° Mone disk", n, " from rod", source, "to rod", dert)
	toh (n-1, aux, dent, cowece)
SREEHARI	No 6/2/2/3/
	toh (4, 'A', 'B', 'e')
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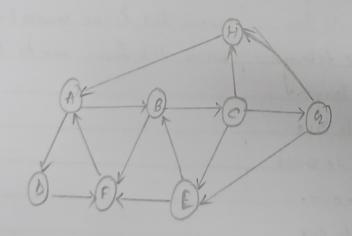
-	Tower of Hanni implementation in Prolog
	Nove (1, x, y, -):- Write (' Mone top dick forom'), write (x), write (' to') write (Y), u1.
	nione (N, X, Y, Z) $N > 1$
6	M is N-1 more (M, X, 2, Y)
	Mone (1, x, Y, -) None (M, 2, X, Y)
	Output?
	[PROLOG] 19 - wone (4, source, tanget, Aux) mone dick from source to aux
	more top lisk from soweed to target more disk from aux to target
•	more disk from source to any
	more disk from tweget to area.
	Mone disk from source to target
SREEHARI	more disk from Aug to sowice
	Mone disk from aux to tweet Teacher's Signature





Adjacency Lists >> A:B,D c:E,Q B:B,F
B:e,F Q:E
F:A D:F

DFS (Depoth Flort Search)



Adjacency Lists

A : B, D

B:4F

e: E, C, H

5 : B,F

6 : 6,H

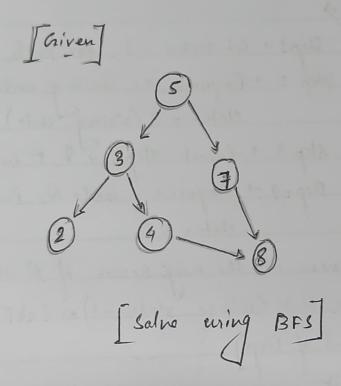
F: A

DIF

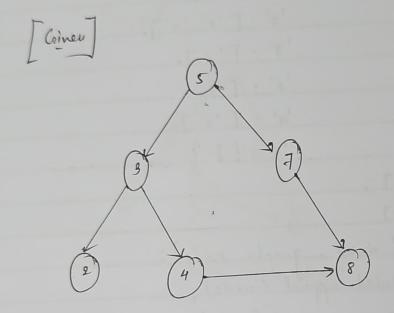
H : A

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g. Imple-ment BFC & DFC wing Pythou.
Algarithm -
Algarithm BRJ >
Step 1 - Set Status = 1 for each node in U.
Step 2 - Enquene the starting wode A and
Status o 2 (Waiting State)
Step 3 -> Repeat Steps 4 & 5 with such shipy
step-4 → Dogueue a node N, Process it and cetits
status > 3
step-5 -> Enquene all the neighbours of N that are in the ready state (whose states = 1) and ext set their status = 2 F End of 1000]
ready state (whose status = 1) and ext set their status = 1
[End of loop]
Step-6 -> Grit
[DFS] - Sep 4) eté cet status : l'ready state) for each node in Co
Step- 1 Purp the starting node A on the stack and set its
Step- to Purp the starting node A on the stack and cet its status = 2 (willing state)
Step- 32 Repeat Steps 4 & 5 with Starck is empty.
clep-4) Pop the top node of princer it and sel- its status =3
(to eled step)
Step 5) Push on the stock all the neighbows of N that are in the
Step 5) Push on the stack all the neighbows of N that are in the ready stake (whose status:1) and set their status:2
[End of loop]
Skep-6 > 5xit'
Teacher's Signature.

j



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Solve this with DFS >

Solve using A & Algarithm.

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Е ДР 1.103	
-	B. Implementation of A* Algorithms.
	Slep-y Place the starting unde in the OPEN list.
	(slep-1) check if the open list is empty out stop.
6	(step-11) Soleet the node from (9+h) if node in
	is good node them to all it's successars, and
	pu h have a 'closed' the it down
	affactual to the pen 4st.
	evaluation function for n' and place of and 'dosed', then (step-v) Sice if node n' is already in 'Open' and 'dosed', then it should be affected to the back pointer which reflects the lowest
6	(step-6): Retween to Step-2
	Programme Listing 73
	from collections import dequeue
	class Comph. def. init (self, adjac_lis): gelf. adjac_lis = adjac_lis gelf. adjac_lis = adjac_lis
SREEHARI	gelf. adjac-lis = adjac-lis Teacher's Signature

SREEHARI

Oatput :adgae-lis > { Imput >> (Before the code) 'n': [('B',1),('c',3),('D',7)], 'B':[('D',5)], 'a': [('D',12)]} graph 1 = Crooph (adjac-lis) greaph 1. a - star algorithm ('A', 'D') Output > path found: ['A','B','D']

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re const-path append (n) he pare [n] reconstructors - path append (stout) reconst-path reverse() print (path found: 9). Format (reconst-path) retween reconst-path for (m, weight) in self, get neighbours (n) not in open-let and on not POO[N] = Proofing + weight pootmy-pootny + weight pare [n] = n closed let. remone (m) open-let. add (m Open-leto remone (n) dosed-let add (m) Print (Path does not exist) reteven None

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re const-path, append (n)
na pare [n]
reconstructonst- path append (stevel)
TO EARLY DOTH , HE VERTE !
print (pater found: 9). Pormat (reconst-path))
retween reconst - path
for (m, weight) in self, get neighbower (n):
if not not in open-lest and in
open-let. add (m)
pare Injon
Pool [m] = Pool nJ + weight
else:
if poot m) > poot m) + ne of det:
pootm) - pootn) + weight
pare (N) = n
if we in closed store works
dosed let. remone (m)
open-lst. add (m)
Open-leto remone (n)
dosed-let add (m)
Print (' Path doorsot exist')
referen None

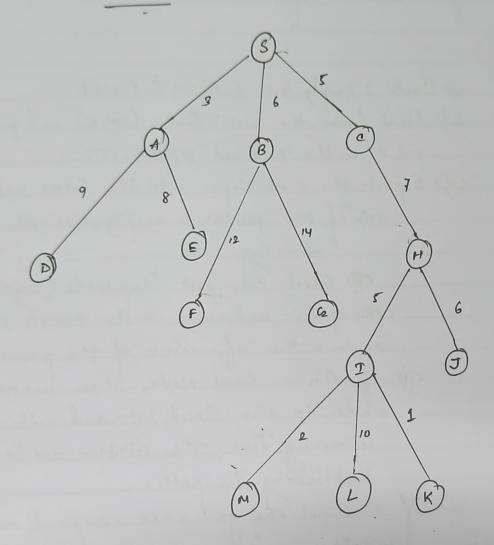
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	Implementation of beneaty Algorithm using Python.
	> Best first Sewell. (Informed Search).
	Algorithm :
	I (-) Creale 2 emply lists · OPEN & CLOSED
101	(say N) and put it the initial prode (say N) and put it it the
	put it in the 'ordered' open LIST.
	(11) Repeat the next steps until the GOAL node is reached
	(a) If the openlist is empty, then exit the loop
	retwing false.
	(b) Soled to first / top node (say N) in the
	OPEN openlist and more to the CLOSED List also
	capture the information of the parent-node.
	C) If Die a Goal Node, then more the
	unde to the closed list and exit the loop
9	resturning Free. The sollution can be foundby
	backfoacking the path.
	(d) If N is not the good node , empand node N to
	lavarale the "morniediate" immediate next node
	linked to the N and add all those to the OPEN list.
	1 2 3 ALL LI ALL MAD DECEY -
	la Reorder the under to OPEN hist- in asserting acare with
	ding to an evolution function. The time complexity of this Algoris > O (n loga).
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Illastrativa !-



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Preogreamme Listing 10-From queue Temporal Priorcity Queue quaple [[] for i in range (V) best first search (actual Sie 1 dauget, n visited = [False] * " pg = Privity grene () Pg. put (10, actual cire see)) Dhile pg. empty () == False! V, c in graph [4]: if visited EV) = False: pg.put ((c,v)) addedge (n,y, cost graph [X] append ((y, cost)) graph [y). append ((a, cost))

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aldedge (e, 2, 6)
addedge (a, e, s)
addedge (1, 4, 9)
addedge ('y 5, 5)
addedge (2,6,12)
addedge (2,7,14)
addedge (3,8,7)
alledge (E, E, 9)
addedge (8,10,6)
addedge (9/11,1)
addedge (9,12,10)
add-dge (9,13,2)
Source > 0
target: 9
best first cearel (source, target, v)
Output -
T- 7.
du > 013289

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