



# Unit – 1

# Introduction to Compiler

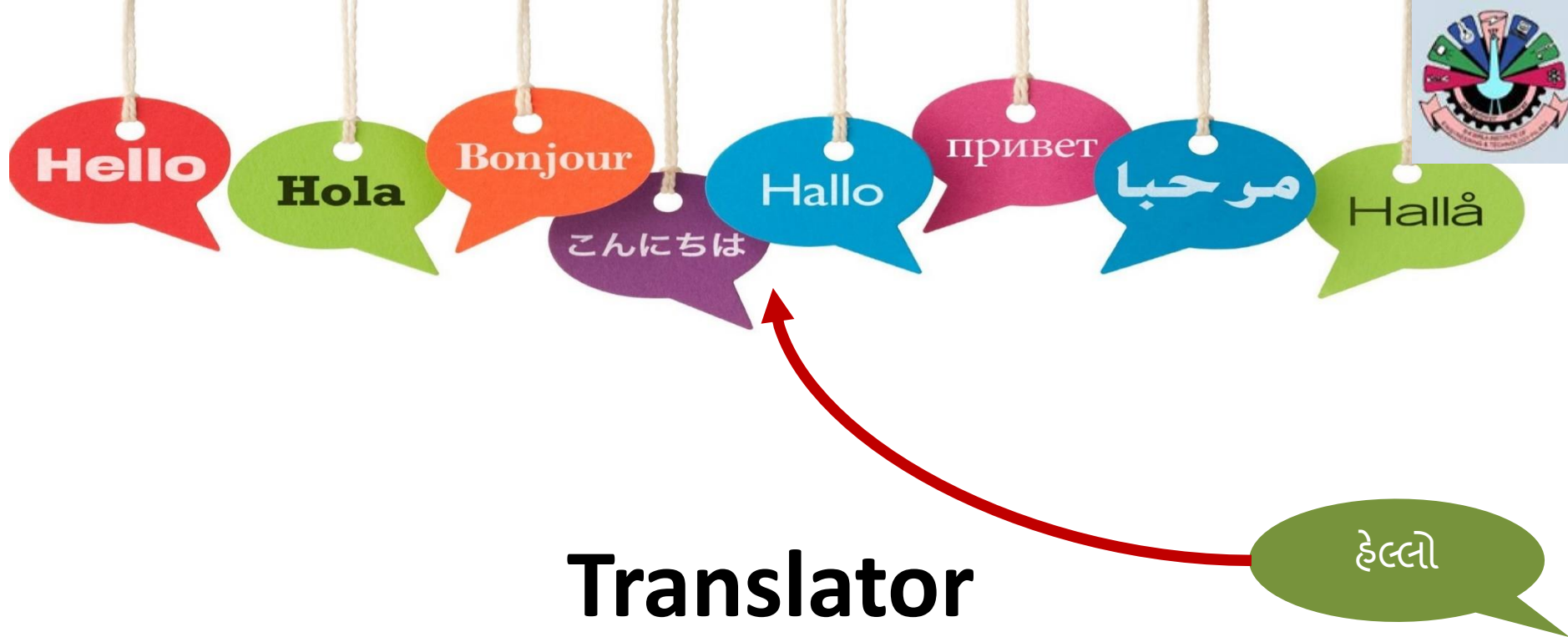
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# Topics to be covered



- Translator
- Analysis synthesis model of compilation
- Phases of compiler
- Difference between compiler & interpreter
- Types of compiler
- Context of compiler (Cousins of compiler)
- Pass structure



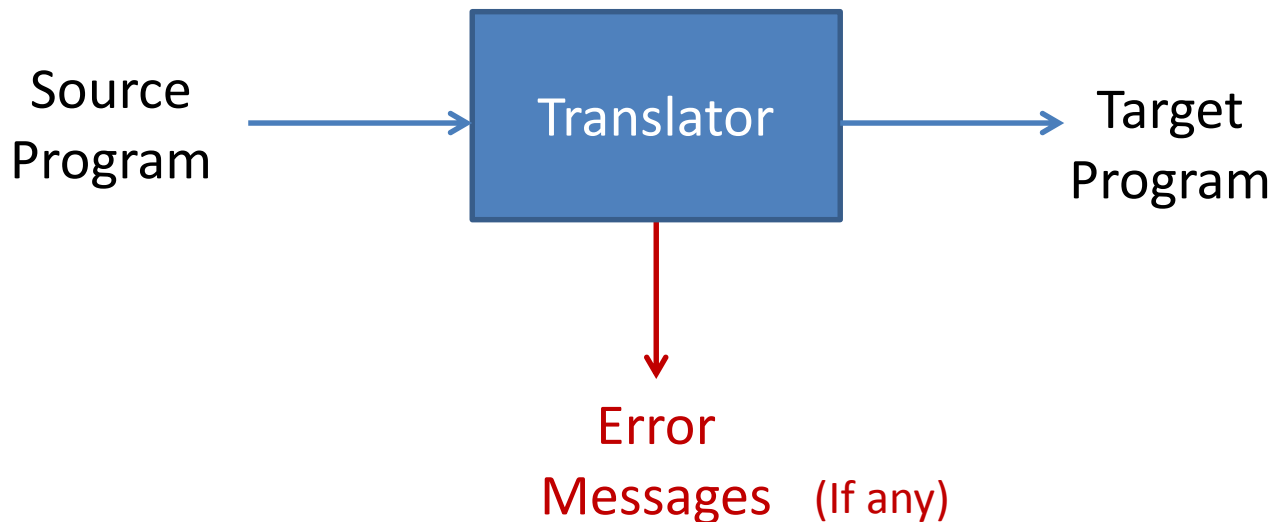
**Translator**

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# Translator



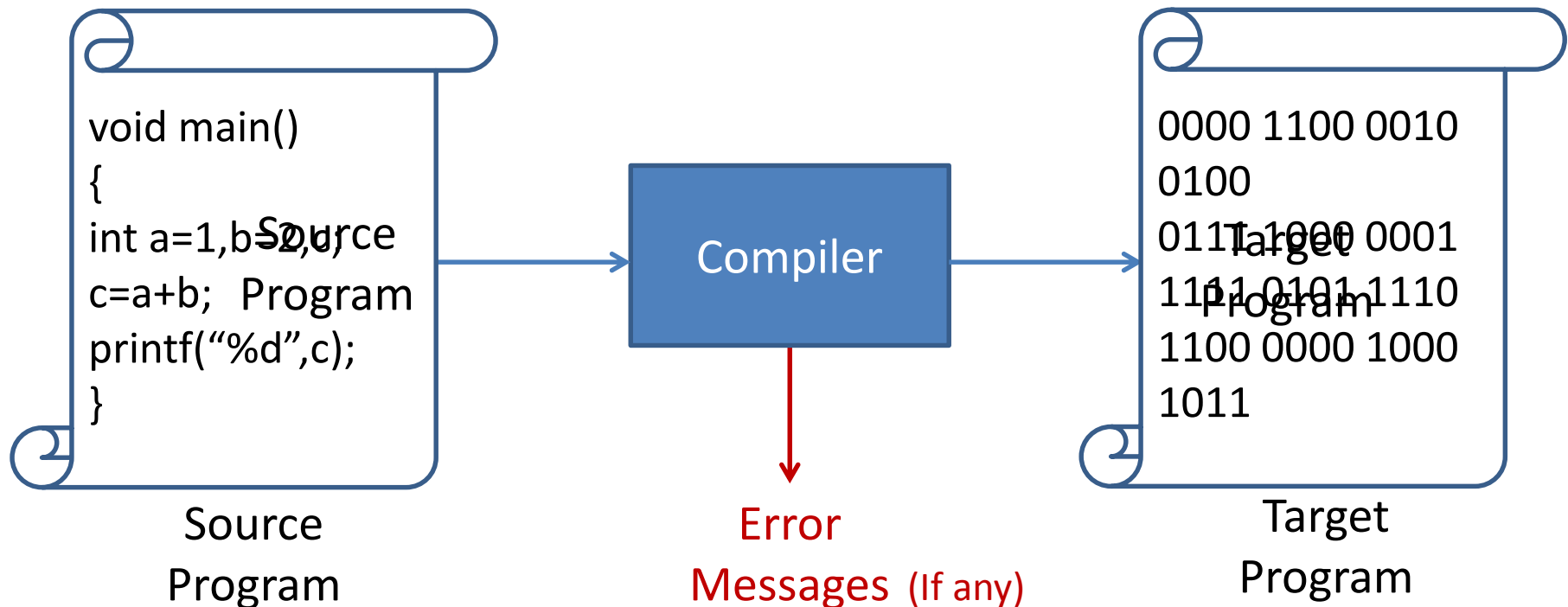
- A translator is a program that takes one form of program as input and converts it into another form.
- Types of translators are:
  1. Compiler
  2. Interpreter
  3. Assembler



# Compiler



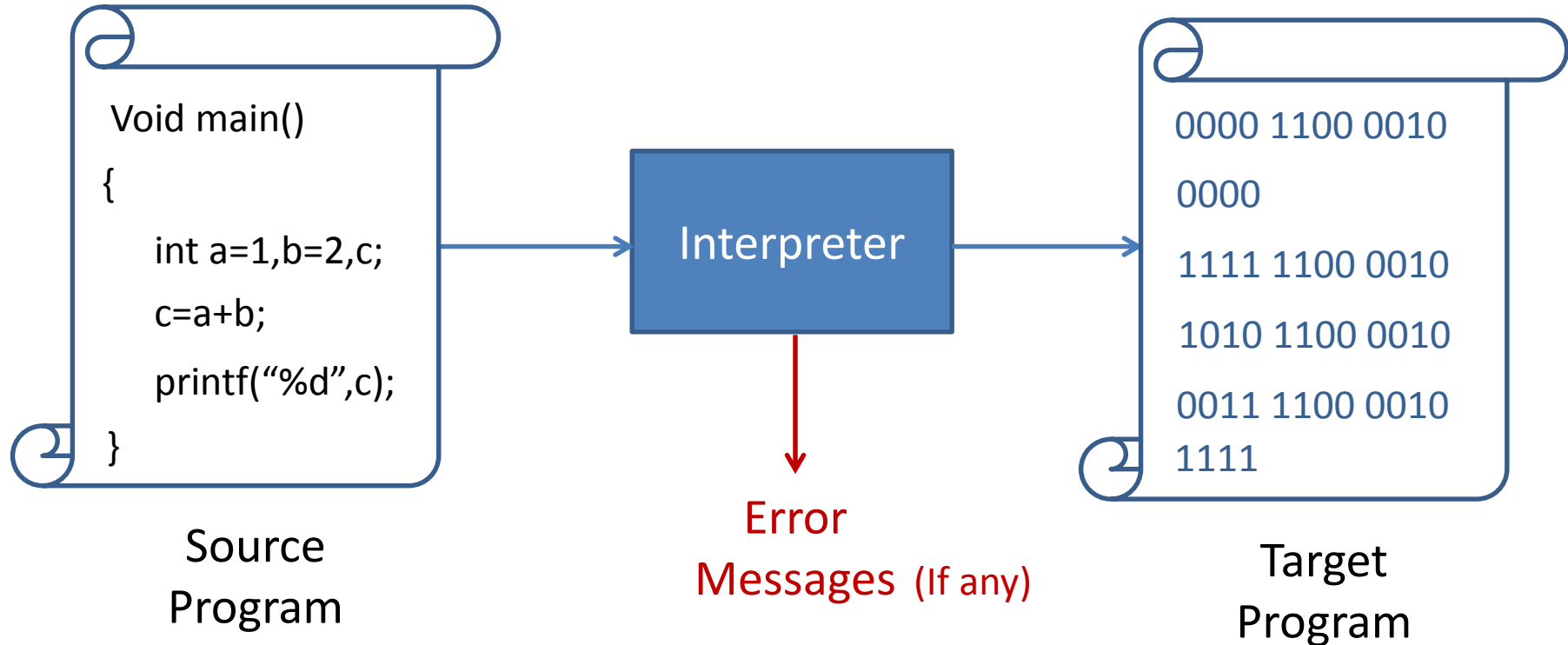
- A compiler is a program that reads a program written in source language and translates it into an equivalent program in target language.



# Interpreter



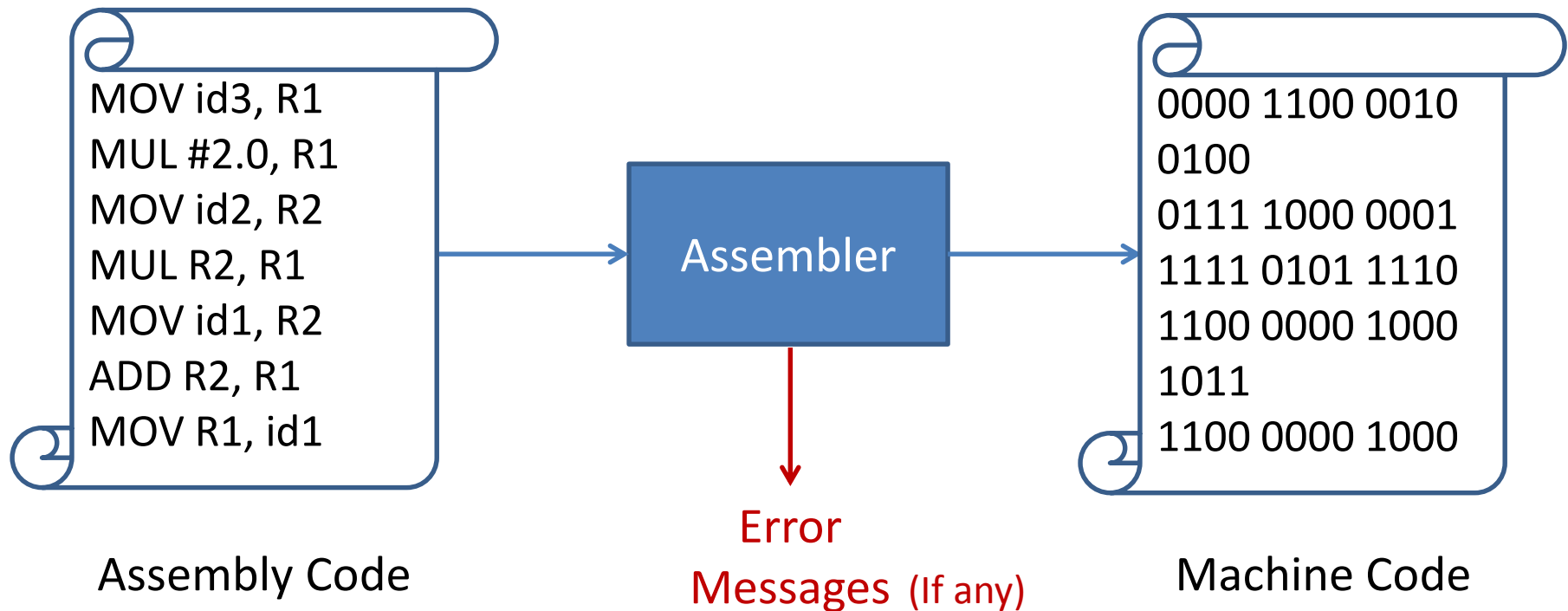
- Interpreter is also program that reads a program written in source language and translates it into an equivalent program in target language line by line.



# Assembler



- Assembler is a translator which takes the assembly code as an input and generates the machine code as an output.





# **Analysis synthesis model of compilation**

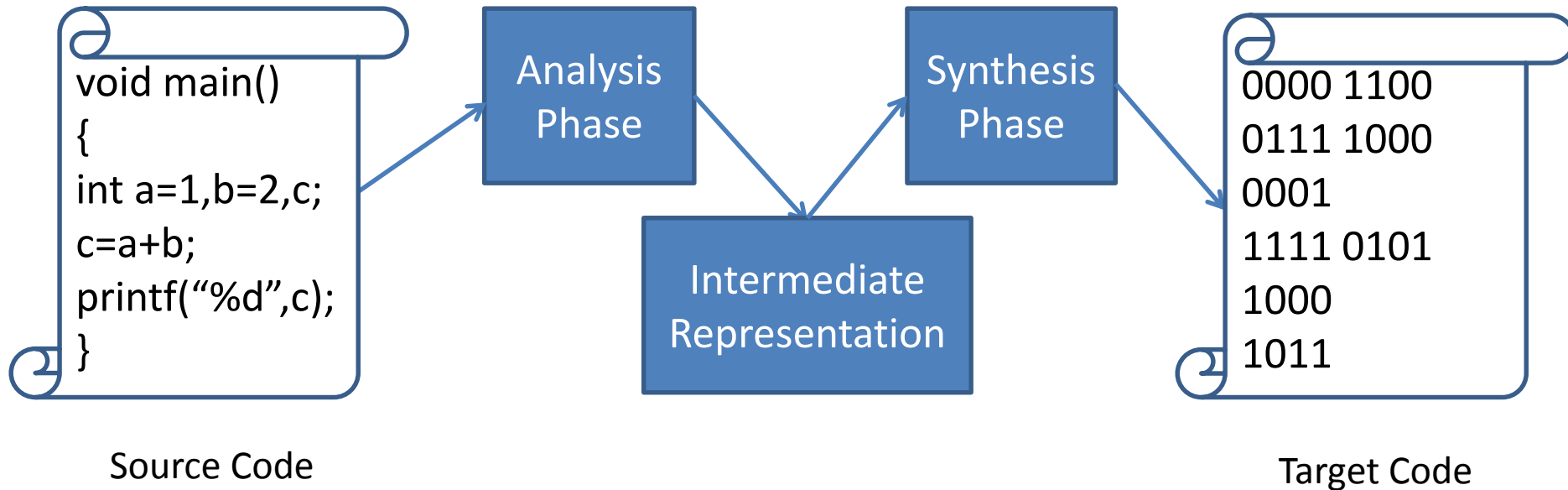


# Analysis synthesis model of compilation



- There are two part of compilation.

1. Analysis Phase
2. Synthesis Phase



# Analysis phase & Synthesis phase



## Analysis Phase

- Analysis part breaks up the source program into constituent pieces and creates an intermediate representation of the source program.
- Analysis phase consist of three sub phases:
  1. Lexical analysis
  2. Syntax analysis
  3. Semantic analysis

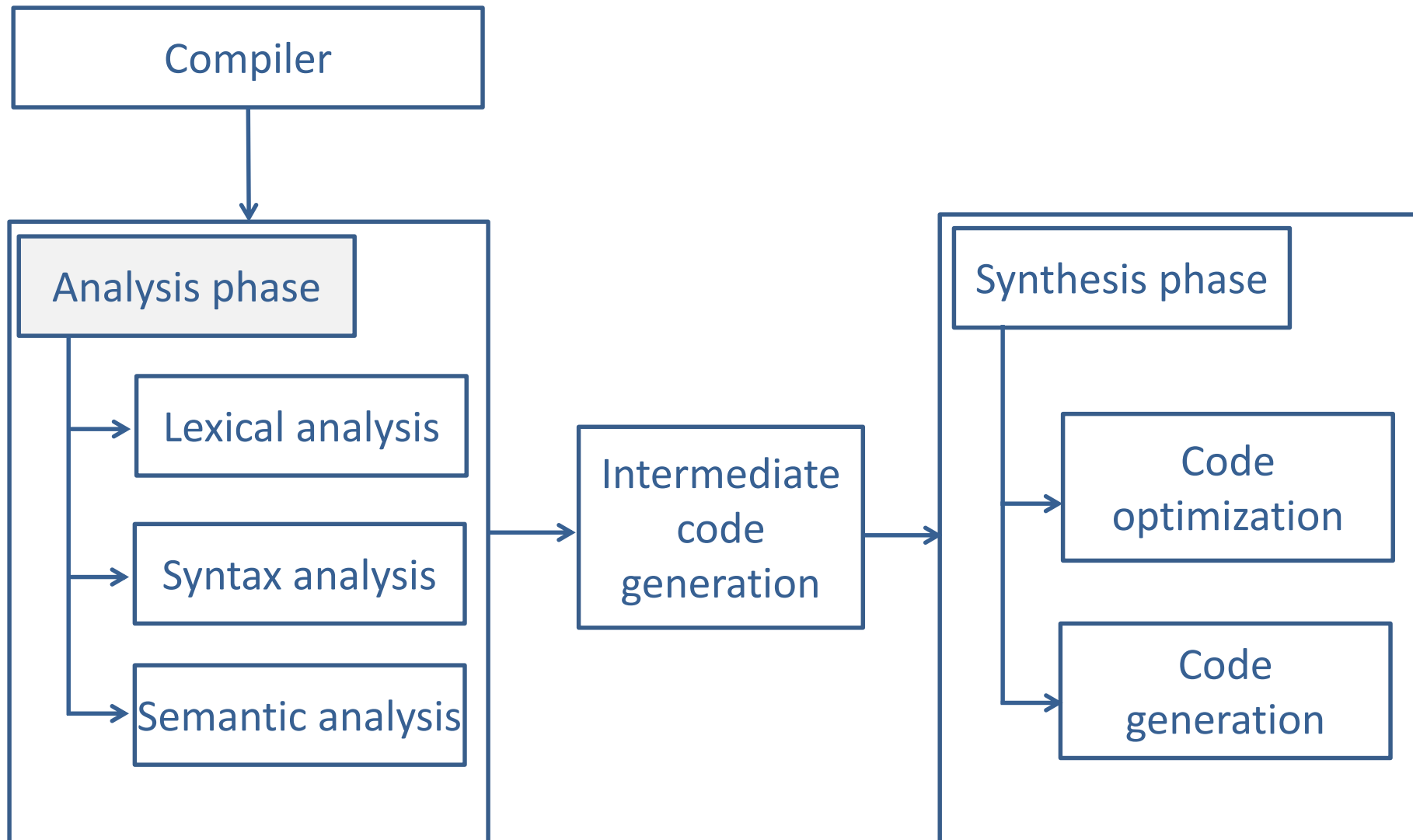
## Synthesis Phase

- The synthesis part constructs the desired target program from the intermediate representation.
- Synthesis phase consist of the following sub phases:
  1. Code optimization
  2. Code generation



# Phases of compiler

# Phases of compiler



# Lexical analysis



- Lexical Analysis is also called *linear analysis* or *scanning*.
- Lexical Analyzer divides the given source statement into the *tokens*.
- Ex: `Position = initial + rate * 60` would be grouped into the following tokens:

`Position` (identifier)

`=` (Assignment symbol)

`initial` (identifier)

`+` (Plus symbol)

`rate` (identifier)

`*` (Multiplication symbol)

`60` (Number)

Position = initial + rate\*60

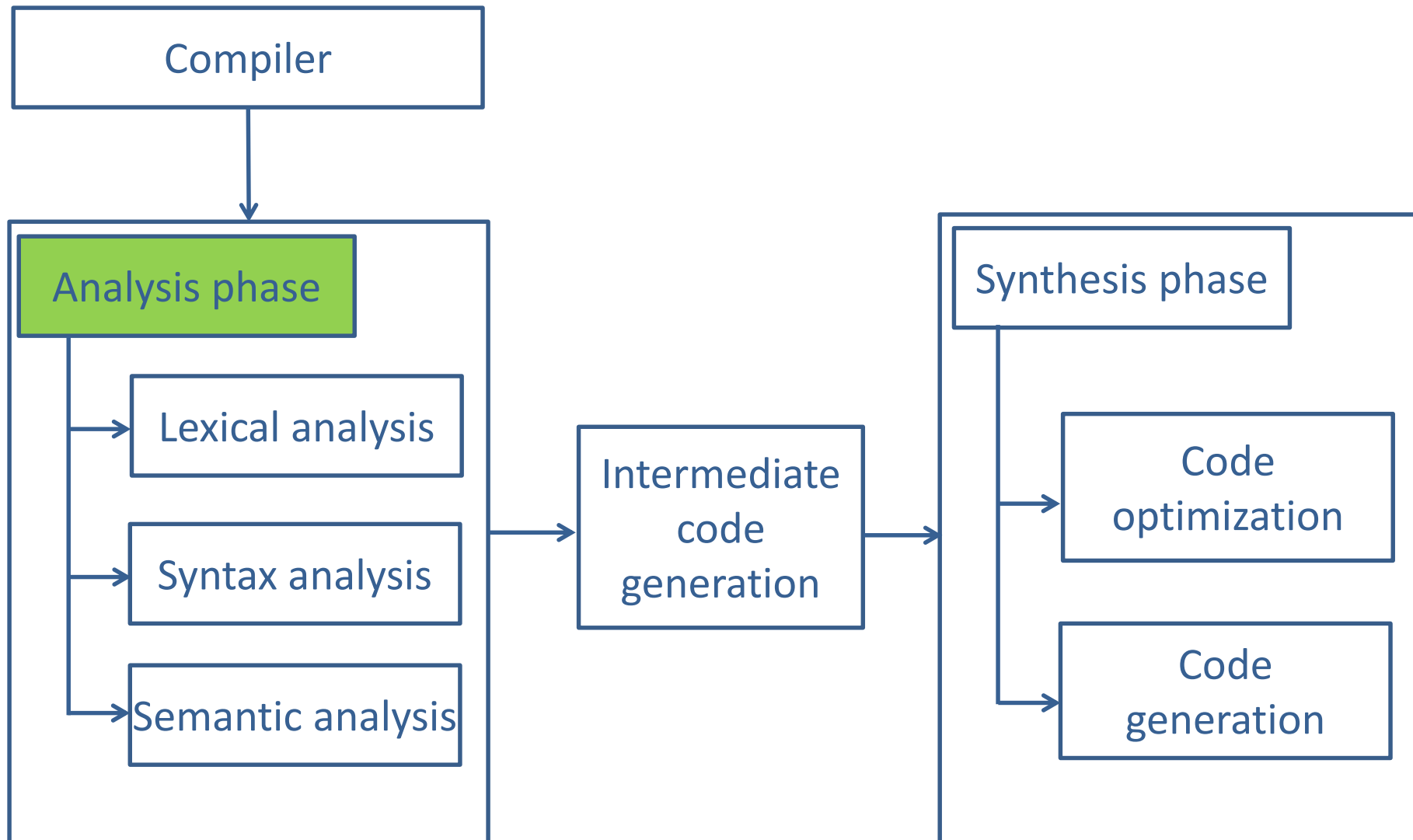


Lexical analysis



`id1 = id2 + id3 * 60`

# Phases of compiler



# Syntax analysis



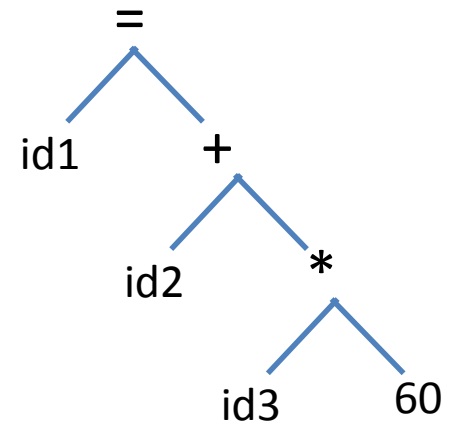
- Syntax Analysis is also called *Parsing* or *Hierarchical Analysis*.
- The syntax analyzer checks each line of the code and spots every tiny mistake.
- If code is error free then syntax analyzer generates the tree.

Position = initial + rate\*60

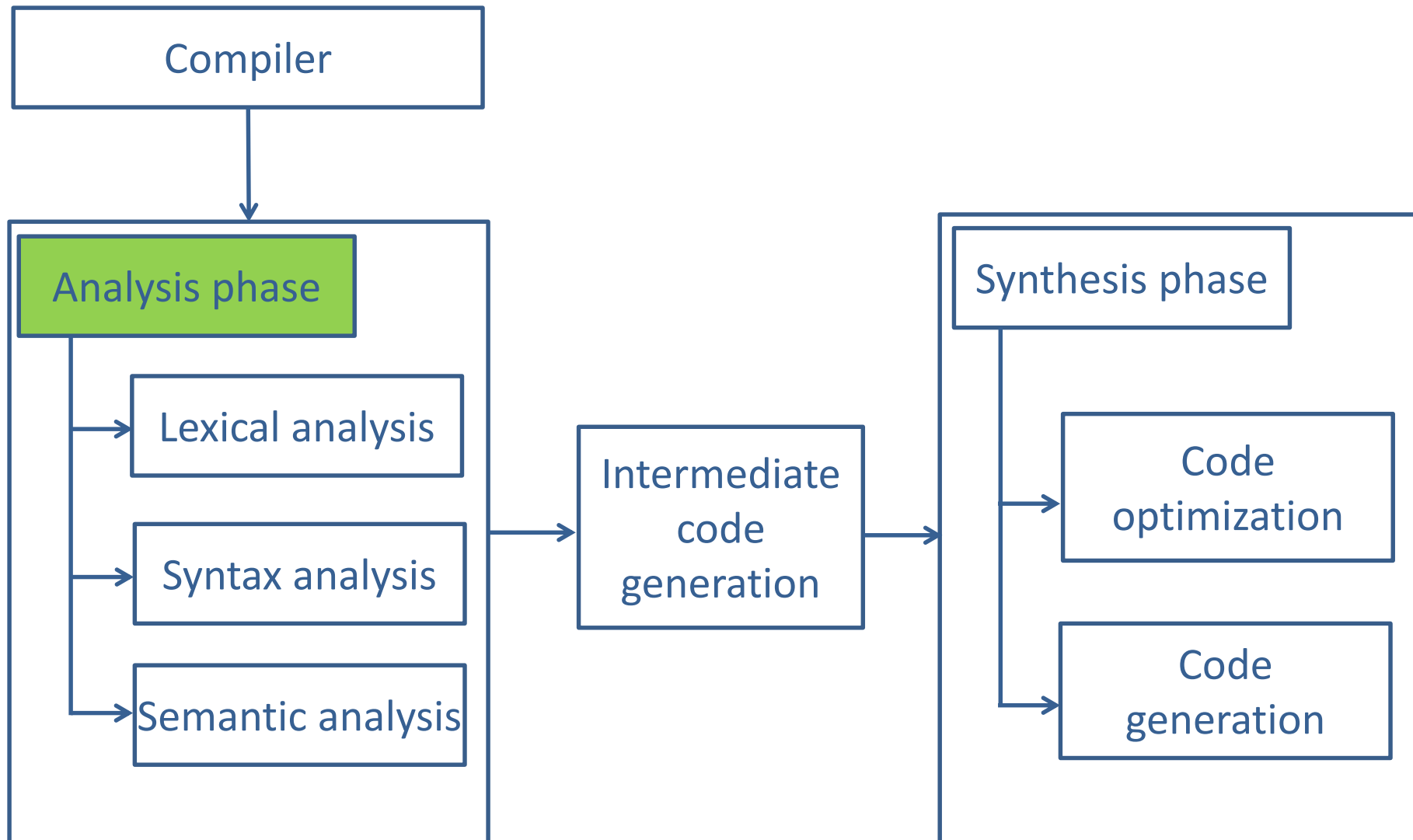
Lexical analysis

id1 = id2 + id3 \* 60

Syntax analysis



# Phases of compiler



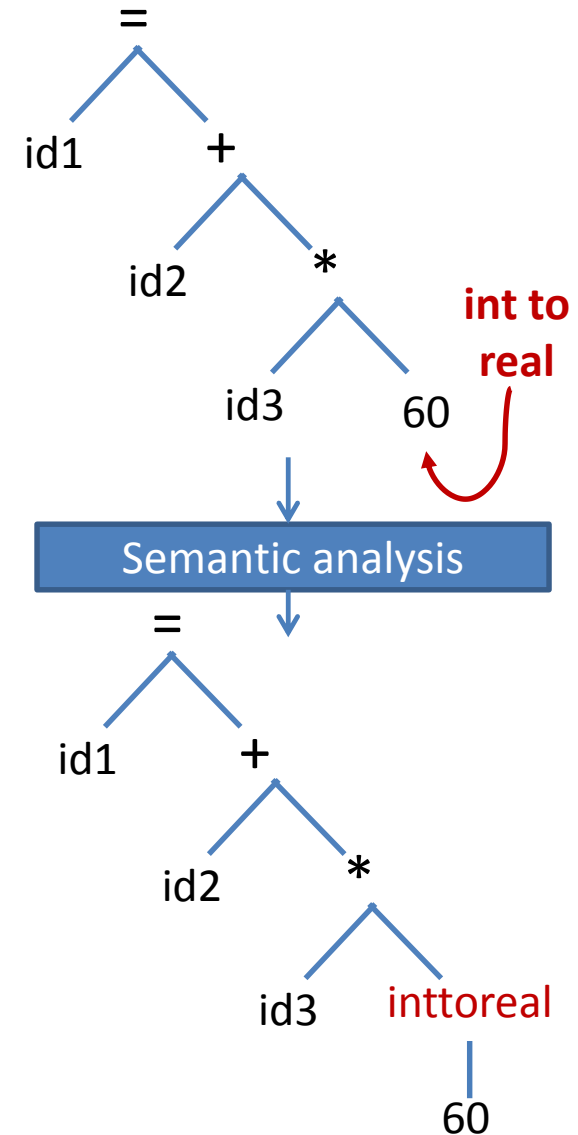


# Semantic analysis

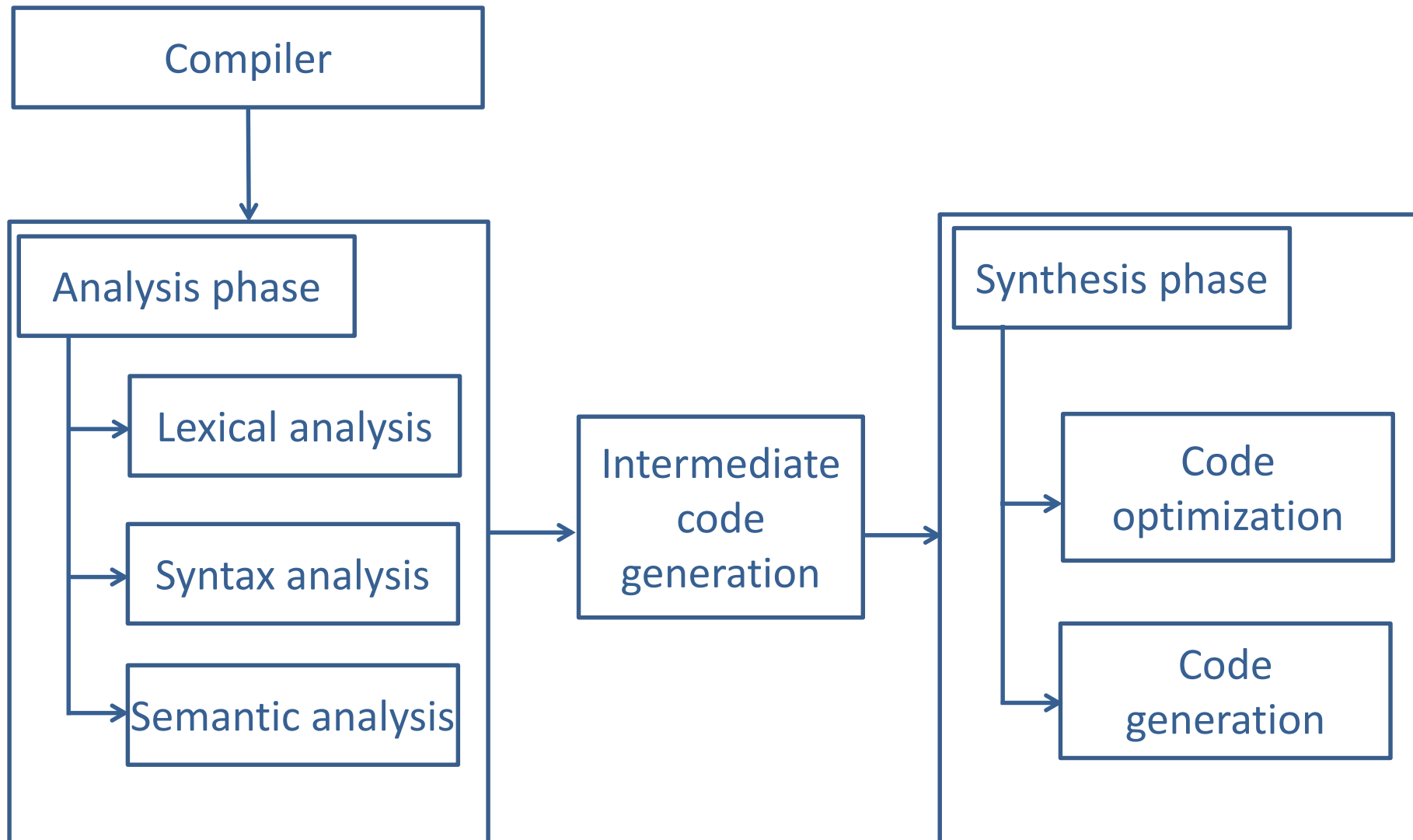


- Semantic analyzer determines the meaning of a source string.
- It performs following operations:
  1. matching of parenthesis in the expression.
  2. Matching of if..else statement.
  3. Performing arithmetic operation that are type compatible.
  4. Checking the scope of operation.

\*Note: Consider id1, id2 and id3 are real



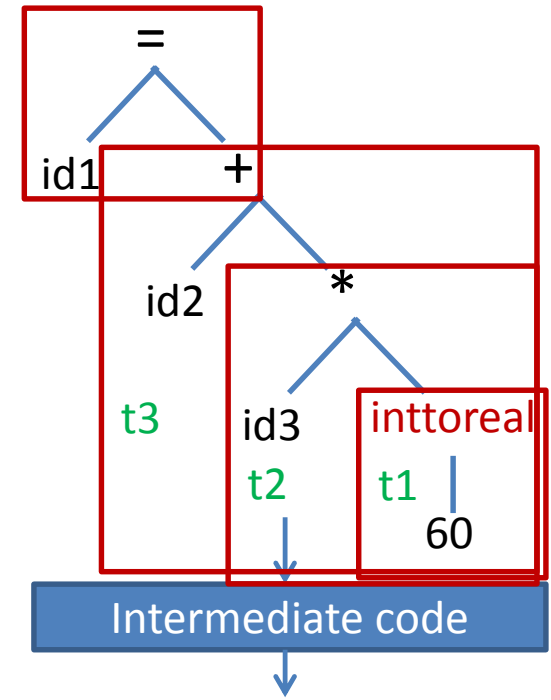
# Phases of compiler



# Intermediate code generator

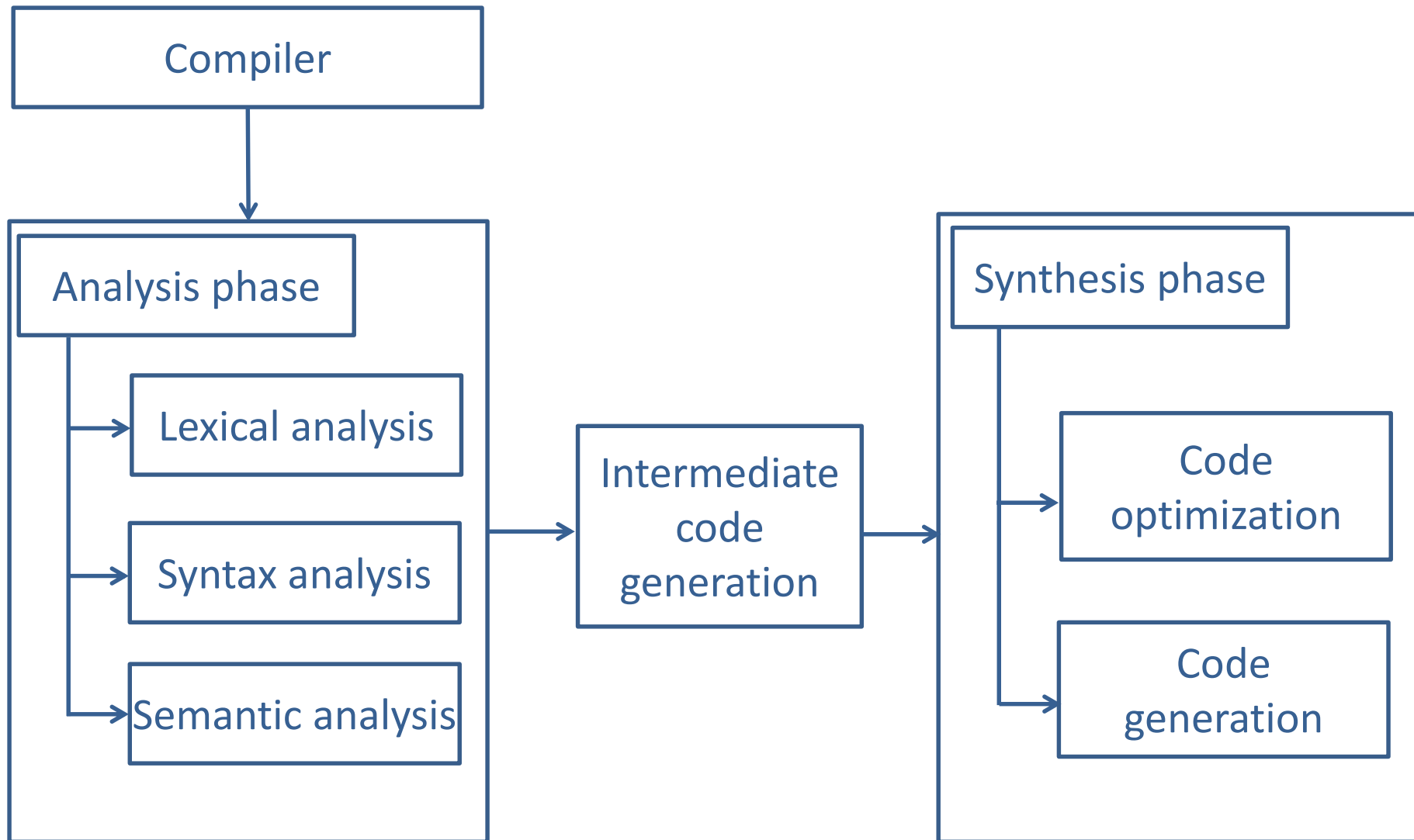


- Two important properties of intermediate code :
  1. It should be easy to produce.
  2. Easy to translate into target program.
- Intermediate form can be represented using “three address code”.
- Three address code consist of a sequence of instruction, each of which has at most three operands.



```
t1= int to real(60)
t2= id3 * t1
t3= t2 + id2
id1= t3
```

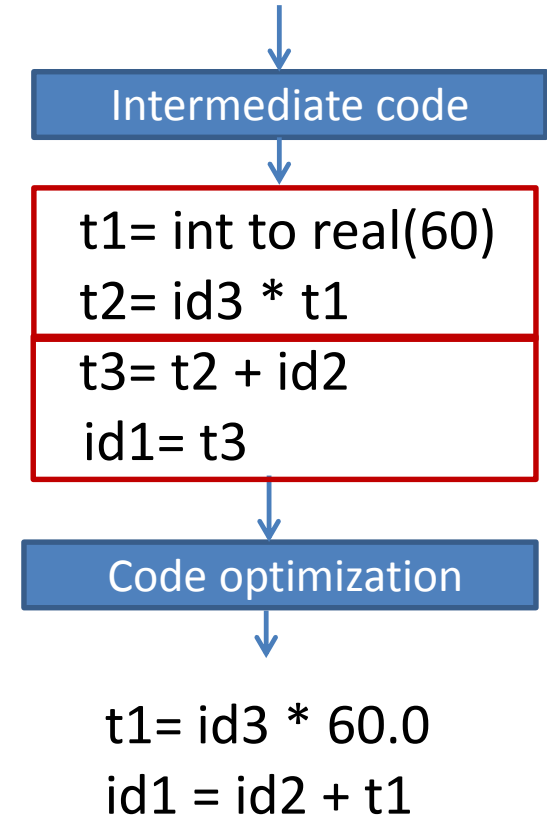
# Phases of compiler



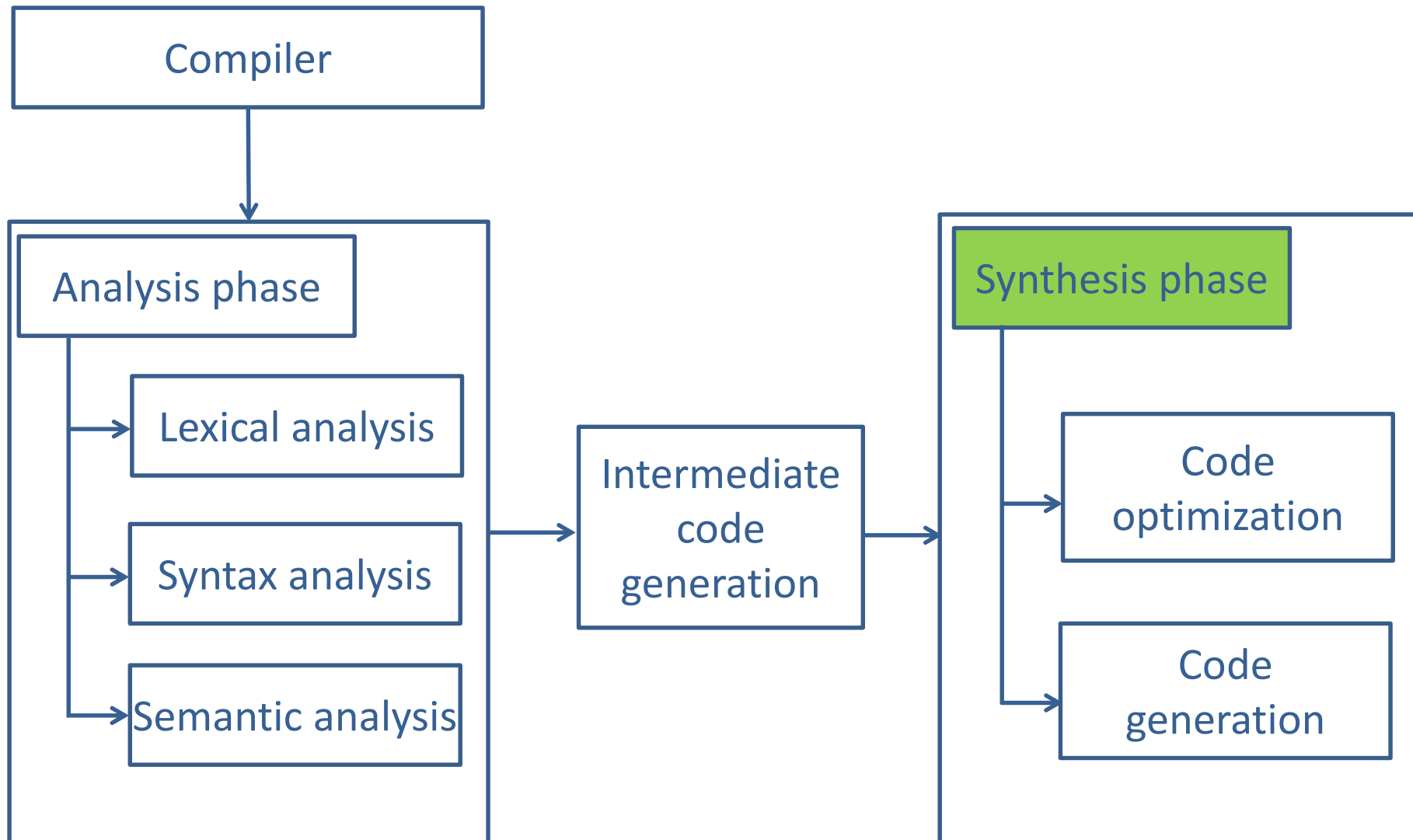
# Code optimization



- It improves the intermediate code.
- This is necessary to have a faster execution of code or less consumption of memory.

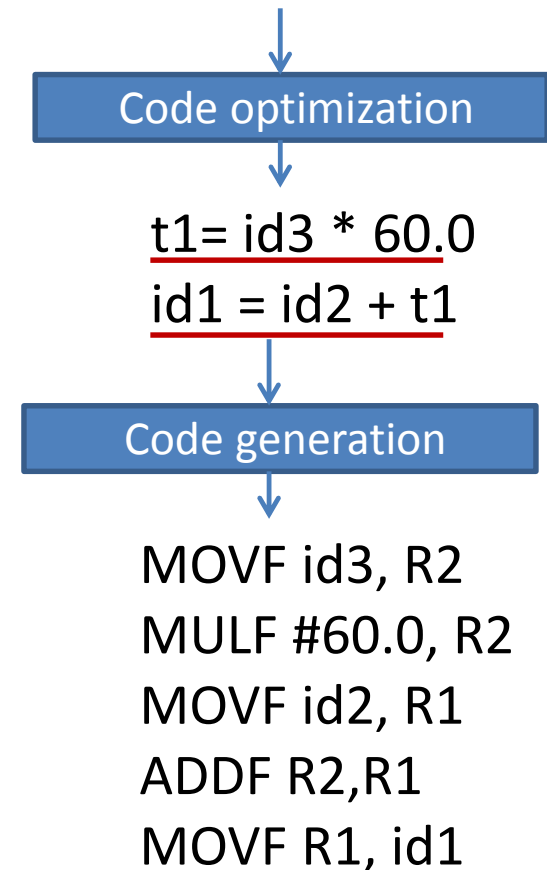


# Phases of compiler

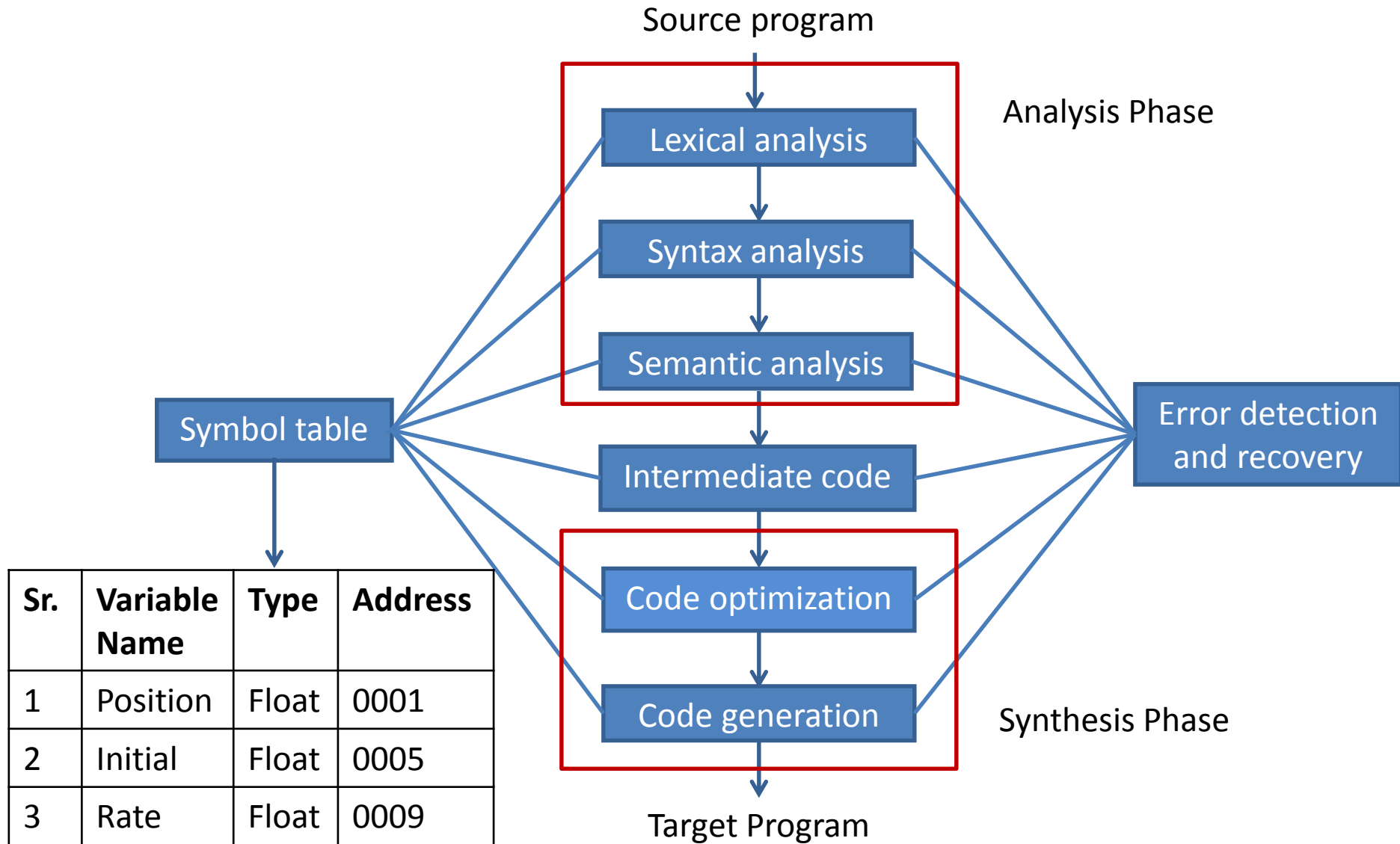


# Code generation

- The intermediate code instructions are translated into sequence of machine instruction.



# Phases of compiler





# Exercise



- Write output of all the phases of compiler for following statements:
  1.  $x = b - c * 2$
  2.  $l = p * n * r / 100$

# Difference between compiler & interpreter



Compiler	Interpreter
Scans the <b>entire program and translates</b> it as a whole into machine code.	It translates program's <b>one statement at a time</b> .
It <b>generates</b> intermediate code.	It <b>does not</b> generate intermediate code.
An error is displayed <b>after entire program</b> is checked.	An error is displayed for <b>every instruction</b> interpreted if any.
Memory requirement is <b>more</b> .	Memory requirement is <b>less</b> .
Example: C compiler	Example: Basic, Python, Ruby

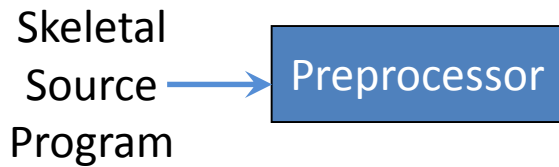
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- The scanning/lexical analysis phase of a compiler performs the task of **reading the source program** as a file of characters and **dividing up into tokens**
  - Usually implemented **as subroutine or co-routine** of parser.
  - Front end of compiler.



# **Context of Compiler (Cousins of compiler)**

# Context of compiler

## (Cousins of compiler)



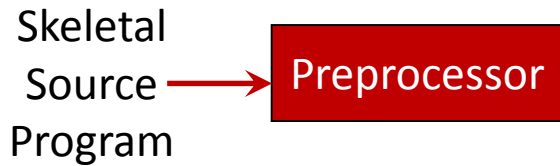
## Preprocessor

It performs the following functions:

1. Macro processing
2. File inclusion
3. Rational preprocessor
4. Language extensions

# Context of compiler

## (Cousins of compiler)



## Preprocessor

1. **Macro processing:** Allows user to **define macros**. Macro is shorthand for longer constructs.

Ex: `#define PI 3.14159265358979323846`

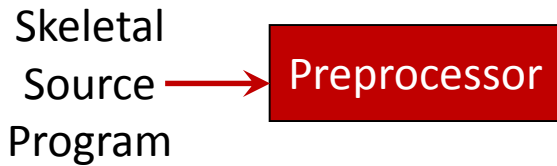
2. **File inclusion:** A preprocessor may **include the header file** into the program.

Ex: `#include<stdio.h>`

3. **Rational preprocessor:** It provides built in macro for construct like **while** statement or **if** statement.

# Context of compiler

## (Cousins of compiler)

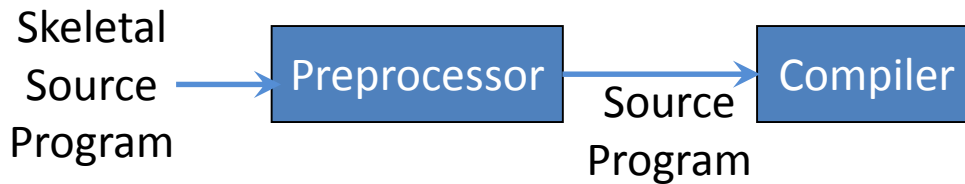


### Preprocessor

4. **Language extensions:** Add capabilities to the language by using built-in macros.
- Ex: the language equal is a database query language embedded in C.
  - Statement beginning with `##` are taken by preprocessor to be database access statement unrelated to C and translated into procedure call on routines that perform the database access.

# Context of compiler

## (Cousins of compiler)



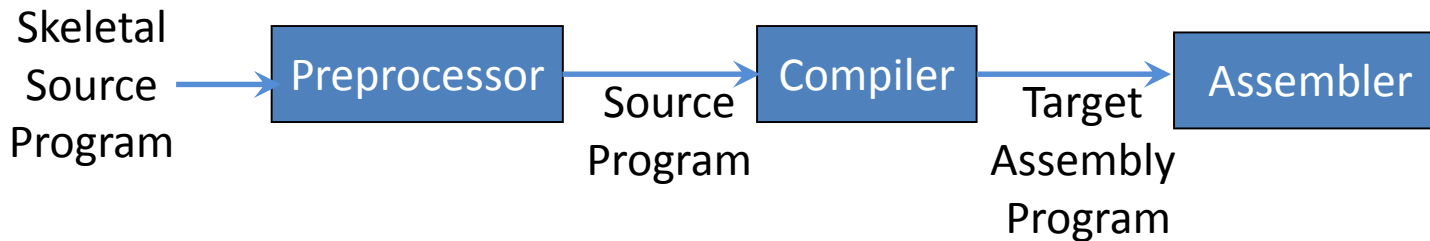
## Compiler

- A compiler is a program that reads a program written in source language and translates it into an equivalent program in target language.



# Context of compiler

## (Cousins of compiler)

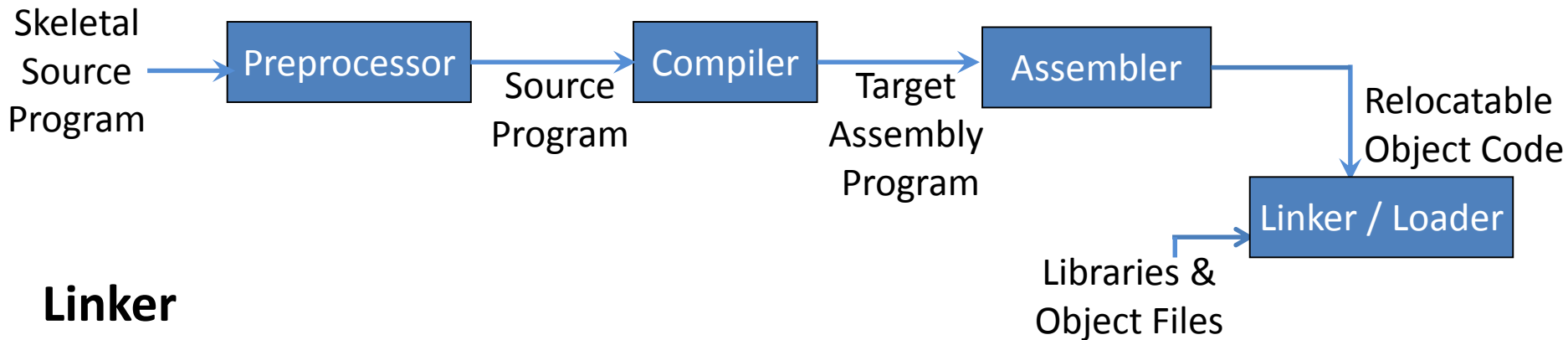


## Assembler

- Assembler is a translator which takes the assembly program (mnemonic) as an input and generates the machine code as an output.

# Context of compiler

## (Cousins of compiler)

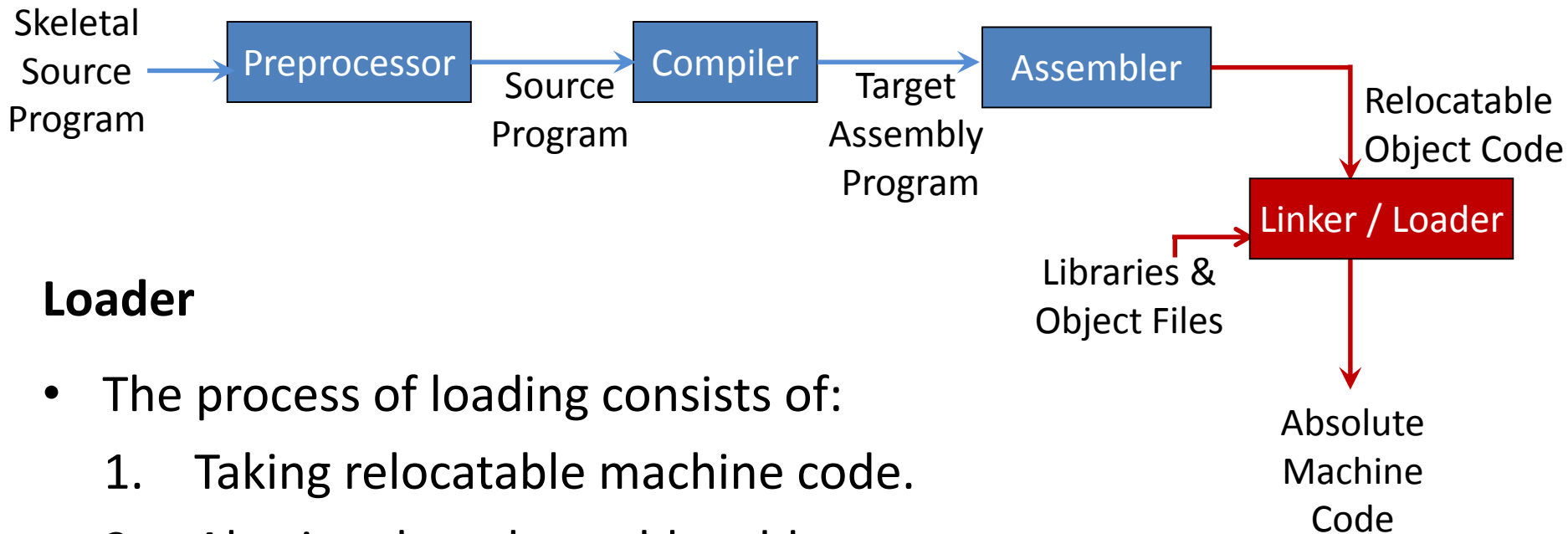


### Linker

- Linker makes a single program from a several files of relocatable machine code.
- These files may have been the result of several different compilation, and one or more library files.

# Context of compiler

## (Cousins of compiler)



### Loader

- The process of loading consists of:
  1. Taking relocatable machine code.
  2. Altering the relocatable address.
  3. Placing the altered instructions and data in memory at the proper location.

# Front end & back end

## (Grouping of phases)



**Front end:** Depends primarily on source language and largely independent of the target machine.

It includes following phases:

1. Lexical analysis
2. Syntax analysis
3. Semantic analysis
4. Intermediate code generation
5. Creation of symbol table

**Back end:** Depends on target machine and do not depend on source program.

It includes following phases:

1. Code optimization
2. Code generation phase
3. Error handling and symbol table operation



# Pass structure

# Pass structure



- One complete scan of a source program is called pass.
- Pass includes **reading an input file** and **writing to the output** file.
- In a single pass compiler analysis of source statement is immediately followed by synthesis of equivalent target statement.
- While in a two pass compiler intermediate code is generated between analysis and synthesis phase.
- It is difficult to compile the source program into single pass due to:  
*forward reference*

# Pass structure



**Forward reference:** A forward reference of a program entity is a **reference to the entity which precedes its definition** in the program.

- This problem can be solved by postponing the generation of target code until more information concerning the entity becomes available.
- It leads to multi pass model of compilation.
- In Pass I: Perform analysis of the source program and note relevant information.
- In Pass II: Generate target code using information noted in pass I.

# Effect of reducing the number of passes



- It is desirable to have a few passes, because it takes time to read and write intermediate file.
- If we group several phases into one pass then memory requirement may be large.





# Types of compiler

# Types of compiler



## 1. One pass compiler

- It is a type of compiler that compiles whole process in one-pass.

## 2. Two pass compiler

- It is a type of compiler that compiles whole process in two-pass.
- It generates intermediate code.

## 3. Incremental compiler

- The compiler which compiles only the changed line from the source code and update the object code.

## 4. Native code compiler

- The compiler used to compile a source code for a same type of platform only.

## 5. Cross compiler

- The compiler used to compile a source code for a different kinds platform.



# End of PART-1(Unit-I)