

Analysis of Algorithms

[5CS4-05/5IT4-05]

Unit 2: Dynamic Programming

Longest Common Subsequence

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Longest Common Subsequence

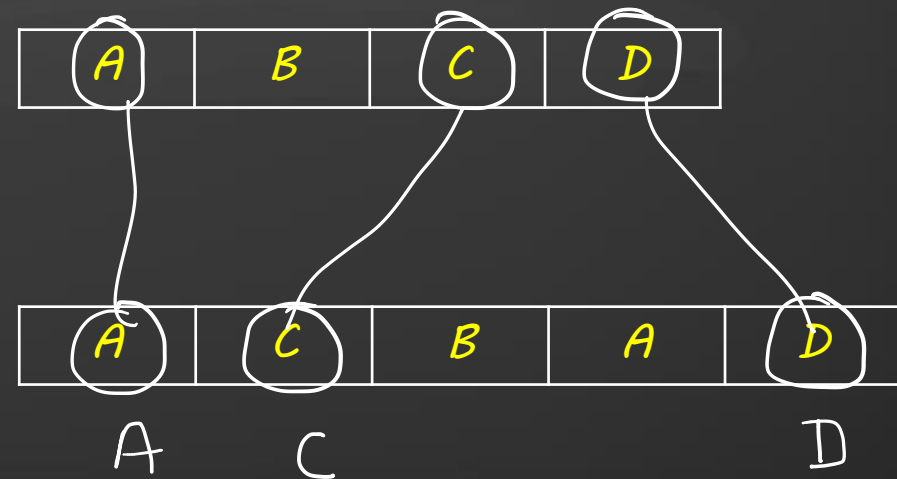
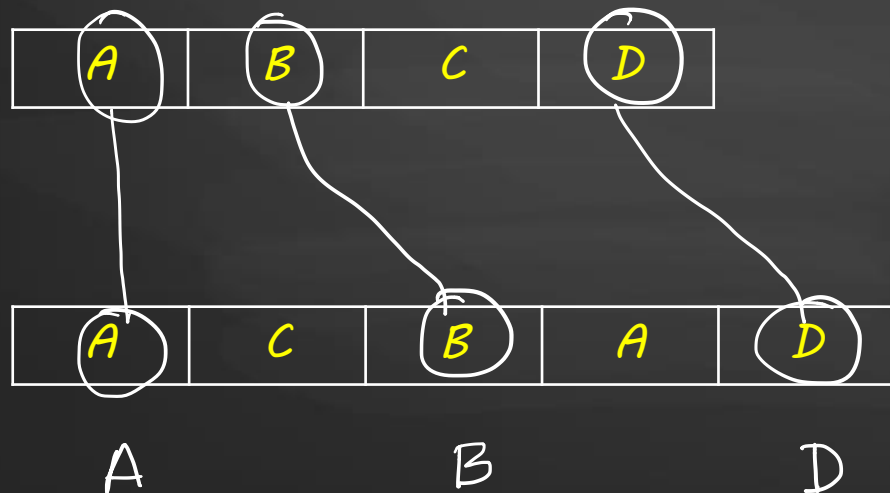
The **longest common subsequence (LCS)** is defined as the longest subsequence that is common to all the given sequences, provided that the elements of the subsequence are not required to occupy consecutive positions within the original sequences.

Longest Common Subsequence

For example, consider the sequences (ABCD) and (ACBAD).

They have 5 length-2 common subsequences: (AB), (AC), (AD), (BD) and (CD)
2 length-3 common subsequences: (ABD) and (ACD)
and no longer common subsequences.

So (ABD) and (ACD) are their longest common subsequences.



① Algorithm to compute LCS length

LCS-LENGTH(X, Y)

```
1   $m = X.length$ 
2   $n = Y.length$ 
3  let  $b[1..m, 1..n]$  and  $c[0..m, 0..n]$  be new tables
4  for  $i = 1$  to  $m$ 
5       $c[i, 0] = 0$ 
6  for  $j = 0$  to  $n$ 
7       $c[0, j] = 0$ 
8  for  $i = 1$  to  $m$ 
9      for  $j = 1$  to  $n$ 
10         if  $x_i == y_j$ 
11              $c[i, j] = c[i - 1, j - 1] + 1$ 
12              $b[i, j] = "\nwarrow"$ 
13         elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
14              $c[i, j] = c[i - 1, j]$ 
15              $b[i, j] = "\uparrow"$ 
16         else  $c[i, j] = c[i, j - 1]$ 
17              $b[i, j] = "\leftarrow"$ 
18  return  $c$  and  $b$ 
```

② Algorithm to print LCS

PRINT-LCS(b, X, i, j)

```
1  if  $i == 0$  or  $j == 0$ 
2      return
3  if  $b[i, j] == \nwarrow$ 
4      PRINT-LCS( $b, X, i - 1, j - 1$ )
5      print  $x_i$ 
6  elseif  $b[i, j] == \uparrow$ 
7      PRINT-LCS( $b, X, i - 1, j$ )
8  else PRINT-LCS( $b, X, i, j - 1$ )
```

Define $c[i, j]$ to be the length of the LCS of $x[1..i]$ and $y[1..j]$

Eg:-

	0	1 A	2 B	3 C	4 D	5 B
0						
1 A						
2 C						
3 B						

```
for  $j = 1$  to  $n$ 
  if  $x_i == y_j$ 
     $c[i, j] = c[i - 1, j - 1] + 1$ 
     $b[i, j] = \nwarrow$ 
  elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
     $c[i, j] = c[i - 1, j]$ 
     $b[i, j] = \uparrow$ 
  else  $c[i, j] = c[i, j - 1]$ 
     $b[i, j] = \leftarrow$ 
```

Define $c[i, j]$ to be the length of the LCS of $x[1..i]$ and $y[1..j]$

	0	1 A	2 B	3 C	4 D	5 B
0	0	0	0	0	0	0
1 A	0	↖ 1	← 1	← 1	← 1	← 1
2 C	0	↑ 1	↑ 1	↖ 2	← 2	← 2
3 B	0	↑ 1	↖ 2	↑ 2	↑ 2	↖ 3

```

for j = 1 to n
  if  $x_i == y_j$ 
     $c[i, j] = c[i - 1, j - 1] + 1$ 
     $b[i, j] = \text{"↖"}$ 
  elseif  $c[i - 1, j] \geq c[i, j - 1]$ 
     $c[i, j] = c[i - 1, j]$ 
     $b[i, j] = \text{"↑"}$ 
  else  $c[i, j] = c[i, j - 1]$ 
     $b[i, j] = \text{"←"}$ 

```

Given the two sequence of characters, $X = \langle A B C B D A B \rangle$, $Y = \langle B D C A B A \rangle$
Find out longest common subsequence

Eg 2:-

		0	1	2	3	4	5	6	7
		\emptyset	A	B	C	B	D	A	B
0	\emptyset								
1	B								
2	D								
3	C								
4	A								
5	B								
6	A								

LCS=

Given the two sequence of characters, $X = \langle A B C B D A B \rangle$, $Y = \langle B D C A B A \rangle$
 Find out longest common subsequence

		0	1	2	3	4	5	6	7
		Ø	A	B	C	B	D	A	B
0	Ø	0	0	0	0	0	0	0	0
1	B	0	↑ 0	↖ 1	← 1	↖ 1	← 1	← 1	↖ 1
2	D	0	↑ 0	↑ 1	↑ 1	↑ 1	↖ 2	← 2	← 2
3	C	0	↑ 0	↑ 1	↖ 2	← 2	↑ 2	↑ 2	↑ 2
4	A	0	↖ 1	↑ 1	↑ 2	↑ 2	↑ 2	↖ 3	← 3
5	B	0	↑ 1	↖ 2	↑ 2	↖ 3	← 3	↑ 3	↖ 4
6	A	0	↖ 1	↑ 2	↑ 2	↑ 3	↑ 3	↖ 4	↑ 4

LCS Length = 4

Given the two sequence of characters, $X = \langle A B C B D A B \rangle$, $Y = \langle B D C A B A \rangle$
 Find out longest common subsequence

		0	1	2	3	4	5	6	7
		Ø	A	B	C	<u>B</u>	<u>D</u>	<u>A</u>	<u>B</u>
0	Ø	0	0	0	0	0	0	0	0
1	B	0	↑ 0	↖ 1	← 1	↖ 1	← 1	← 1	↖ 1
2	D	0	↑ 0	↑ 1	↑ 1	↑ 1	↖ 2	← 2	← 2
3	C	0	↑ 0	↑ 1	↖ 2	← 2	↑ 2	↑ 2	↑ 2
4	A	0	↖ 1	↑ 1	↑ 2	↑ 2	↑ 2	↖ 3	← 3
5	B	0	↑ 1	↖ 2	↑ 2	↖ 3	← 3	↑ 3	↖ 4
6	A	0	↖ 1	↑ 2	↑ 2	↑ 3	↑ 3	↖ 4	↑ 4

LCS = B D A B

Queries ?