



PROGRAMMING LANGUAGE



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Language Overview

The language needs to define the data types of variables being used.

The language supports various arithmetic operation and operators like + , - , / , * and ^ .These operators have their own precedence as in general mathematics.

The basic rules for identifiers are as follows:

- 1 .The first character must be alphabet.
- 2 .All other characters can be either numbers, alphabets or underscore (_).
- 3.LiPi is case sensitive.

The language has associativity l to r for operators of same precedence i.e. left to right.

1 .Basic Data Types:

LiPi supports various data types such as:

1. int for storing integer
2. char for storing characters
3. point for storing floating points

Which is similar to data types of the C language.

2 .Arithmetic Operations:

The language supports various arithmetic operation and operators like + , - , / , * and ^ .These operators have their own precedence as in general mathematics.

Operator	Operation	Example
+	Addition	a+b
-	Subtraction	x-y
*	Multiplication	a1*b1
%	Modulo(Remainder)	a_1*q_1
/	Divide	a/b
^	Power(exponent)	x^g

Operators with same precedence have left to right associativity.

The precedence and associativity of operators is given below.

Operator	Precedence	Associativity
+, -	1	L → R
*, / , %	2	L→R
^	3	R →L

3.Keywords

There are many keywords used in LiPi and are used to perform various tasks and declarations.

List of a few keywords and their use:

Keyword	Purpose
fx	For declaration of a function
main	For definition of the main function(to be executed first)Which controls the flow of program
int	For declaration of a integer variable
point	For declaration of a floating or a rational number
char	For declaration of a character variable
let	For assignments
NULL	Representing void(void data type as in C)
show	to display to the standard console

Grammar:

BNF Versions of Expression Grammar

```
<assign> → <id> = <expr>
<expr> → <expr> + <item>
        | <expr> - <item>
        | <item>
<item> → <item> * <element>
        | <item> / <element>
        | <element>
<element> → <exp> ^ <element>
           | <exp>
<exp> → (<expr>)
       | <id>
<id> → <letter> |
      <id> | <letter> | <digit> | <underscore>
<letter> → A | B | C | D | ..... | Z | a | b | c | d | ..... | z
<digit> → 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9
<underscore> → _
```

EBNF Versions of Expression Grammar

```
<expr> → <item> { (+ | -) <item> }
<item> → <element> { (* | /) <element> }
<element> → <exp> { ^ <exp> }
<exp> → (<expr>)
       | <id>
```

$\langle id \rangle \rightarrow \langle letter \rangle \{ \langle letter \rangle \mid \langle digits \rangle \mid \langle underscore \rangle \}$
 $\langle letter \rangle \rightarrow A \mid B \mid C \mid D \mid \dots \mid Z \mid a \mid b \mid c \mid d \mid \dots \mid z$
 $\langle digit \rangle \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$
 $\langle underscore \rangle \rightarrow _$

A sample program of the language :

```

fx main(NULL)
{
    int a,b,c,d;
    let a=b+c*d;
    show (a);
}

```

Sample statement for the Grammar:

Let us take an arithmetic expression: $A = B + C * D$

$$A = B + C * D$$

Derivation of statement from the Grammar:

$\langle assign \rangle \rightarrow \langle id \rangle = \langle expr \rangle$
 $A = \langle expr \rangle$
 $A = \langle expr \rangle + \langle item \rangle$
 $A = \langle item \rangle + \langle item \rangle$
 $A = \langle element \rangle + \langle item \rangle$
 $A = \langle exp \rangle + \langle item \rangle$
 $A = \langle id \rangle + \langle item \rangle$
 $A = B + \langle item \rangle$
 $A = B + \langle item \rangle * \langle element \rangle$

$A = B + \langle \text{element} \rangle * \langle \text{element} \rangle$

$A = B + \langle \text{exp} \rangle * \langle \text{element} \rangle$

$A = B + \langle \text{id} \rangle * \langle \text{element} \rangle$

$A = B + C * \langle \text{element} \rangle$

$A = B + C * \langle \text{exp} \rangle$

$A = B + C * \langle \text{id} \rangle$

$A = B + C * D$