



# Red-Black Tree- Deletion

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# Red Black Tree- Deletion

- To perform operation, we first execute the deletion algorithm for binary search trees
- Thus, the node which is deleted is the parent of an external node.
- If node is red, it won't violate any property
- If node is a leaf, it won't violate any property
- Otherwise, if node is black and has a child, it will violate property 2, 3, and 4
- For property 2, set the color of root to black after deletion

# Red Black Tree- Deletion

- To fix property 3 and 4:
  - From now on, let's call the deleted node to be  $z$
  - If  $z$ 's child  $x$  (which is the replacing node) is red, set  $x$  to black. Done!
  - If  $x$  is black, add another black to  $x$ , so that  $x$  will be a doubly black node, and property 3 and 4 are fixed. But property 1 is violated

# Red Black Tree- Deletion

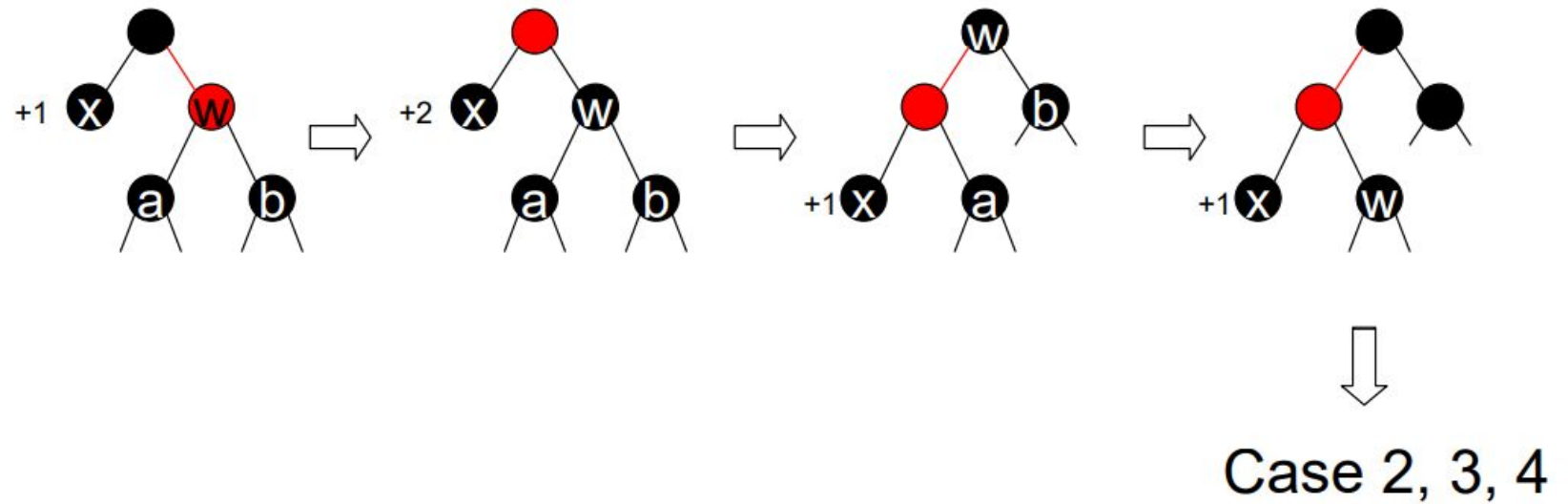
- To fix property 1, we will consider if
  - $x$  is a left child or right child
  - The color of  $x$ 's sibling  $w$  is red or black
  - The colors of  $w$ 's children
- We consider  $x$  is a left child first, the other case can be done by symmetric operation

# Red Black Tree- Deletion

- There are 4 cases:
  - Case 1: w is red
  - Case 2: w is black, both w's children are black
  - Case 3: w is black, w's left child is red, w's right child is black
  - Case 4: w is black, w's right child is red

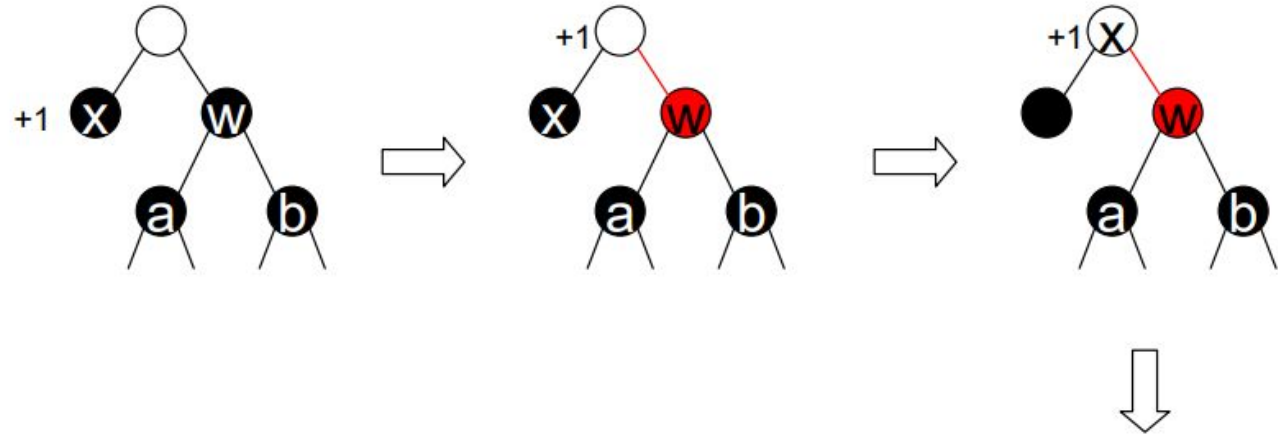
# Red Black Tree - Deletion

- Case 1: w is red



# Red Black Tree - Deletion

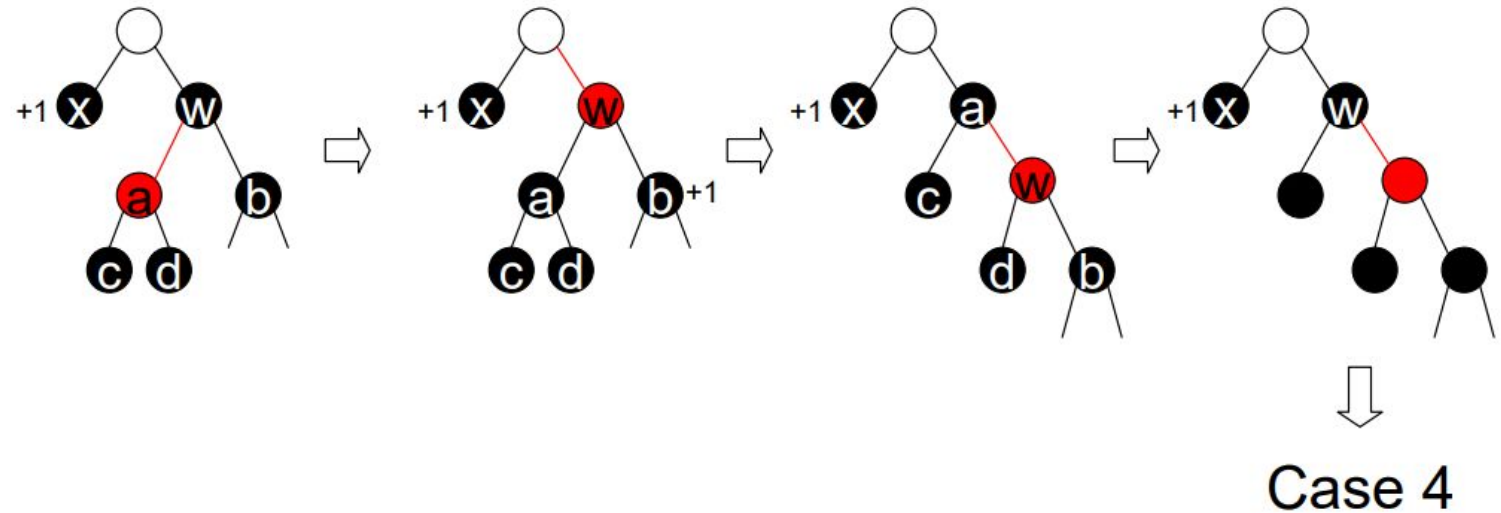
- Case 2: w is black, both w's children are black



Recursively delete x

# Red Black Tree - Deletion

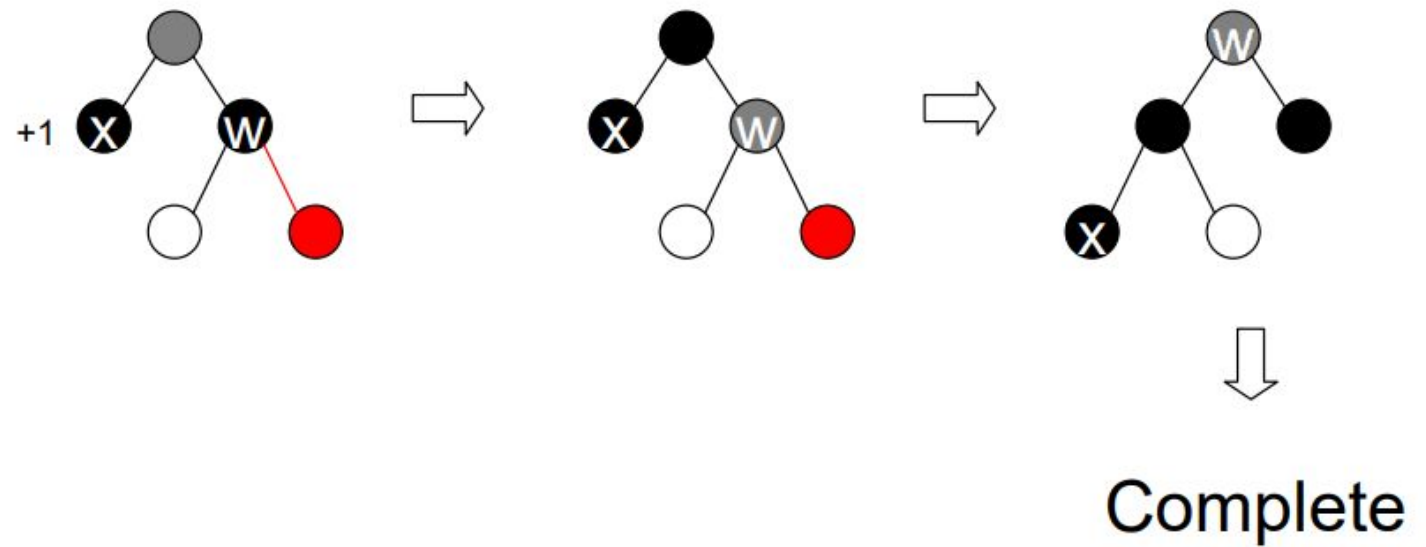
- Case 3: w is black, w's left child is red, w's right child is black





# Red Black Tree - Deletion

- Case 4: w is black, w's right child is red



# Red-Black Tree Deletion - Pseudocode

- RB-Delete( $T, z$ )
  1. if  $z \rightarrow \text{left} = \text{null}$  or  $z \rightarrow \text{right} = \text{null}$
  2. then  $y \leftarrow z$
  3. else  $y \leftarrow \text{TREE-SUCCESSOR}(z)$
  4. if  $y \rightarrow \text{left} \neq \text{null}$
  5. then  $x \leftarrow y \rightarrow \text{left}$
  6. else  $x \leftarrow y \rightarrow \text{right}$
  7.  $x \rightarrow p \leftarrow y \rightarrow p$
  8. if  $y \rightarrow p = \text{null}$
  9. then  $T \rightarrow \text{root} \leftarrow x$
  10. else if  $y = y \rightarrow p \rightarrow \text{left}$
  11. then  $y \rightarrow p \rightarrow \text{left} \leftarrow x$
  12. else  $y \rightarrow p \rightarrow \text{right} \leftarrow x$
  13. if  $y \neq z$
  14. then  $z \rightarrow \text{key} \leftarrow y \rightarrow \text{key}$
  15. copy  $y$ 's data into  $z$
  16. if  $y \rightarrow \text{color} = \text{BLACK}$
  17. then **RB-DELETE-FIXUP**( $T, x$ )
  18. return  $y$

# Red-Black Tree Deletion - Pseudocode

- **RB-DELETE-FIXUP(T, x)**

1. while  $x \neq T \rightarrow \text{root}$  and  $x \rightarrow \text{color} = \text{BLACK}$
2.     do if  $x = x \rightarrow p \rightarrow \text{left}$
3.         then  $w \leftarrow x \rightarrow p \rightarrow \text{right}$
4.     if  $w \rightarrow \text{color} = \text{RED}$
5.         then  $w \rightarrow \text{color} \leftarrow \text{BLACK}$  Case 1
6.      $x \rightarrow p \rightarrow \text{color} \leftarrow \text{RED}$  Case 1
7.     LEFT-ROTATE(T,  $x \rightarrow p$ ) Case 1
8.      $w \leftarrow x \rightarrow p \rightarrow \text{right}$  Case 1
9.     if  $w \rightarrow \text{left} \rightarrow \text{color} = \text{BLACK}$  and  
        $w \rightarrow \text{right} \rightarrow \text{color} = \text{BLACK}$
10.         then  $w \rightarrow \text{color} \leftarrow \text{RED}$  Case 2
11.      $x \leftarrow x \rightarrow p$  Case 2
12.     else if  $w \rightarrow \text{right} \rightarrow \text{color} = \text{BLACK}$
13.         then  $w \rightarrow \text{left} \rightarrow \text{color} \leftarrow \text{BLACK}$  Case 3
14.      $w \rightarrow \text{color} \leftarrow \text{RED}$  Case 3
15.     RIGHT-ROTATE(T,  $w$ ) Case 3
16.      $w \leftarrow x \rightarrow p \rightarrow \text{right}$  Case 3
17.      $w \rightarrow \text{color} \leftarrow x \rightarrow p \rightarrow \text{color}$  Case 4
18.      $x \rightarrow p \rightarrow \text{color} \leftarrow \text{BLACK}$  Case 4
19.      $w \rightarrow \text{right} \rightarrow \text{color} \leftarrow \text{BLACK}$  Case 4
20.     LEFT-ROTATE(T,  $x \rightarrow p$ ) Case 4
21.      $x \leftarrow T \rightarrow \text{root}$  Case 4
22.     else (same as then clause with "right" and  
       "left" exchanged)
23.          $x \rightarrow \text{color} \leftarrow \text{BLACK}$

# Red Black Tree – Deletion Summary

In all cases, except 2, deletion can be completed by a simple rotation/ recoloring.

In case 2, the height of the subtree reduces and so we need to proceed up the tree

- If we proceed up the tree, we only need to recolor/rotate.

Complexity-  $O(\log n)$

# Red-Black Trees to 2-4 Trees

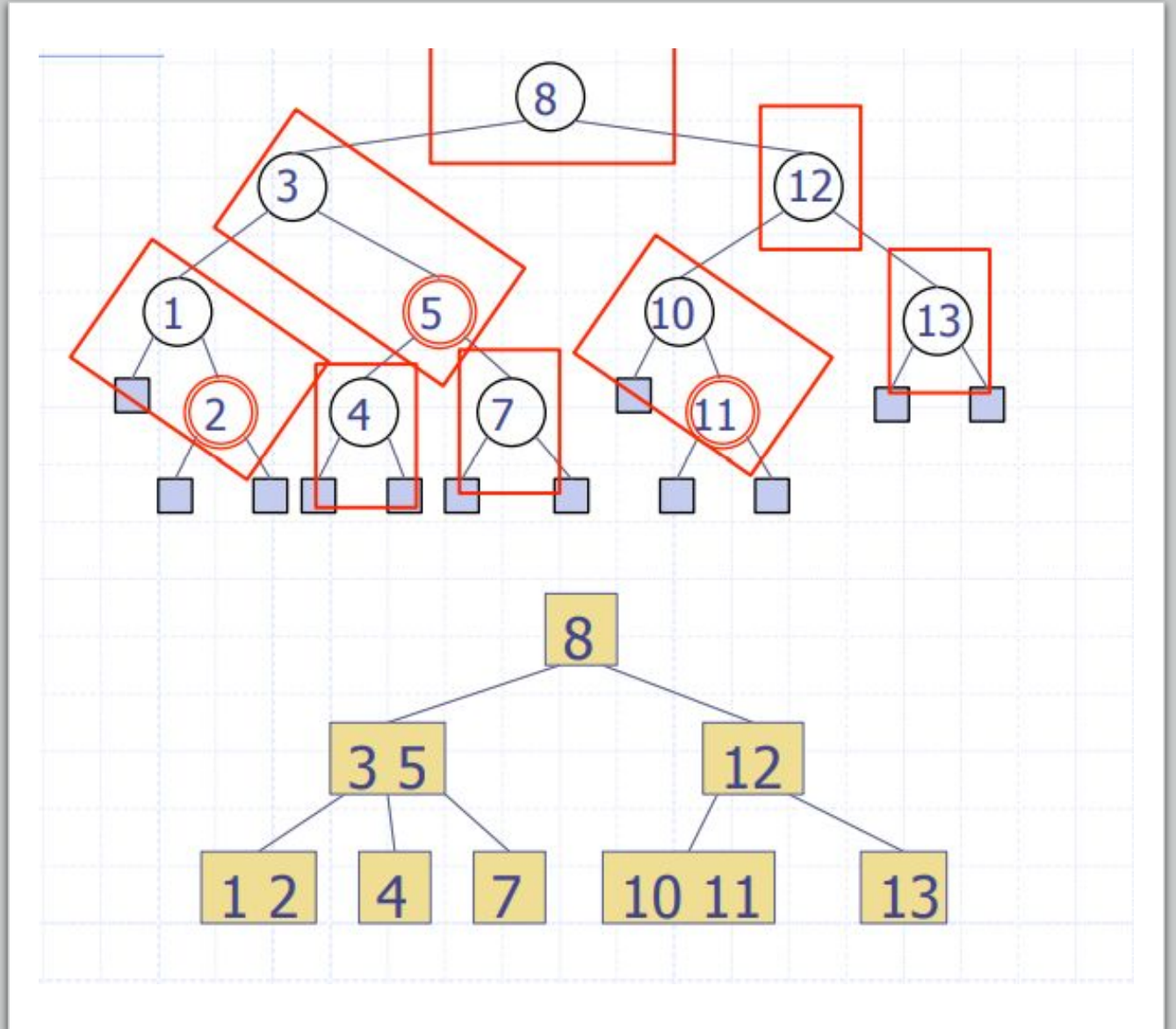
Any red-black tree can be converted into a 2-4 tree

Take a black node and its red children (at most 2) and combine them into one node of a 2-4 tree.

Each node thus formed has at least 1 and at most 3 keys

Since black depth of all external nodes is the same, in the resulting 2-4 tree all the external nodes will be at the same level.

# Red-Black Tree to 2-4 Tree - Example



# 2-4 Trees to Red-Black Trees

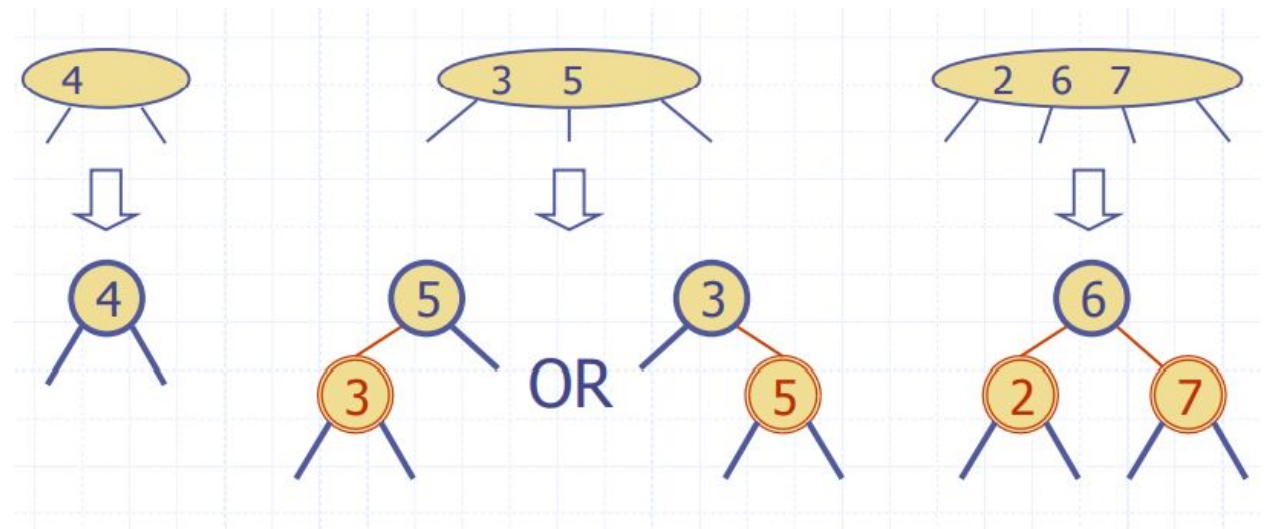
Any 2-4 tree can be converted into a red-black tree

We replace a node of the 2-4 tree with one black node and 0/1/2 red nodes which are children of the black node.

The height of 2-4 tree is the black depth of the red-black tree created.

Every red node has a black child.

# 2-4 Tree to Red-Black Trees





# 2-4 Trees to Red-Black Trees - Example

