

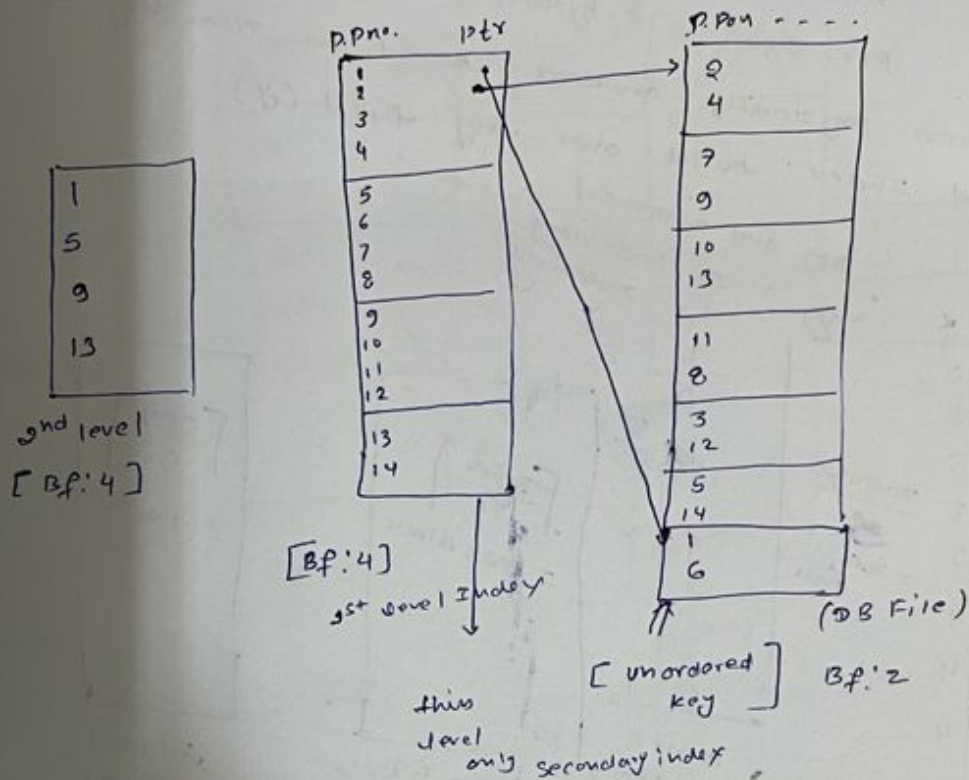
* Secondary Index.

① Search key used for index may not ordered field [un-ordered field] and

(10/29/2021)

lec-31

② Search key may be key (or) Non-key.



- * Secondary index always dense index.
- * more than one secondary index possible for DB Table.
- * ~~more~~ Case I/O cost :- $(k+1)$ block.
worst
- * I/O cost of secondary index more than I/O cost of primary index.

Q.

of records 30000

Block size : 1024 Bytes

Record size : 100 Bytes

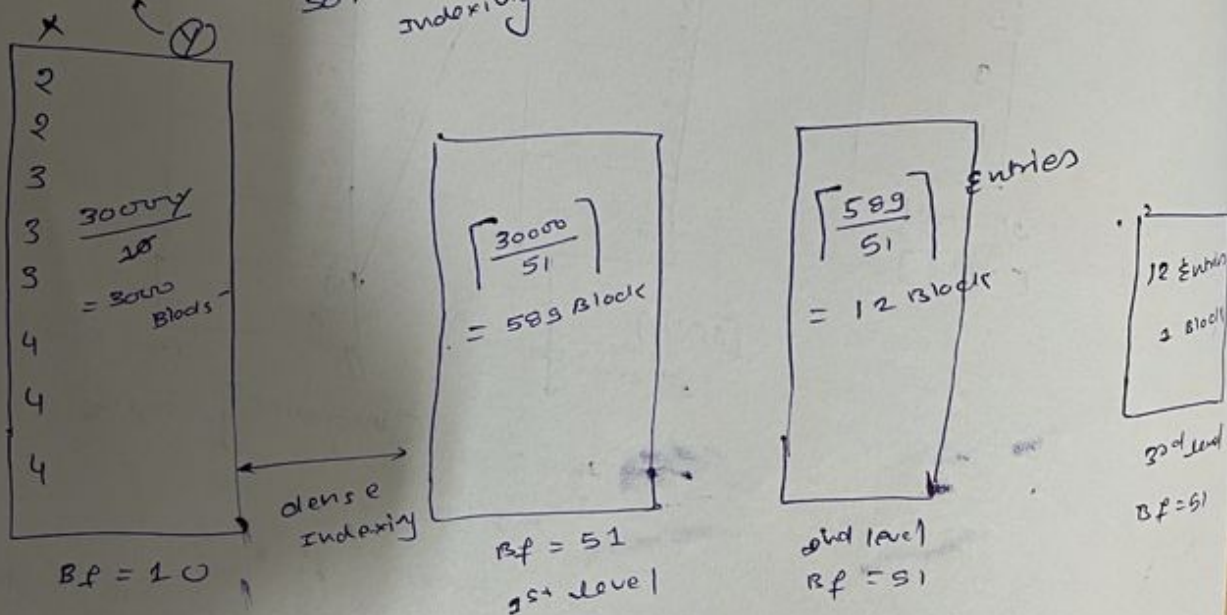
Search key : 12 Bytes.

Pts. size : 8 bytes.

[unspanned org. Used] Limit

Records physically ordered based on non-key field (X)
and Index build over key field (Y),

Solution
key and unordered
so, secondary indexing



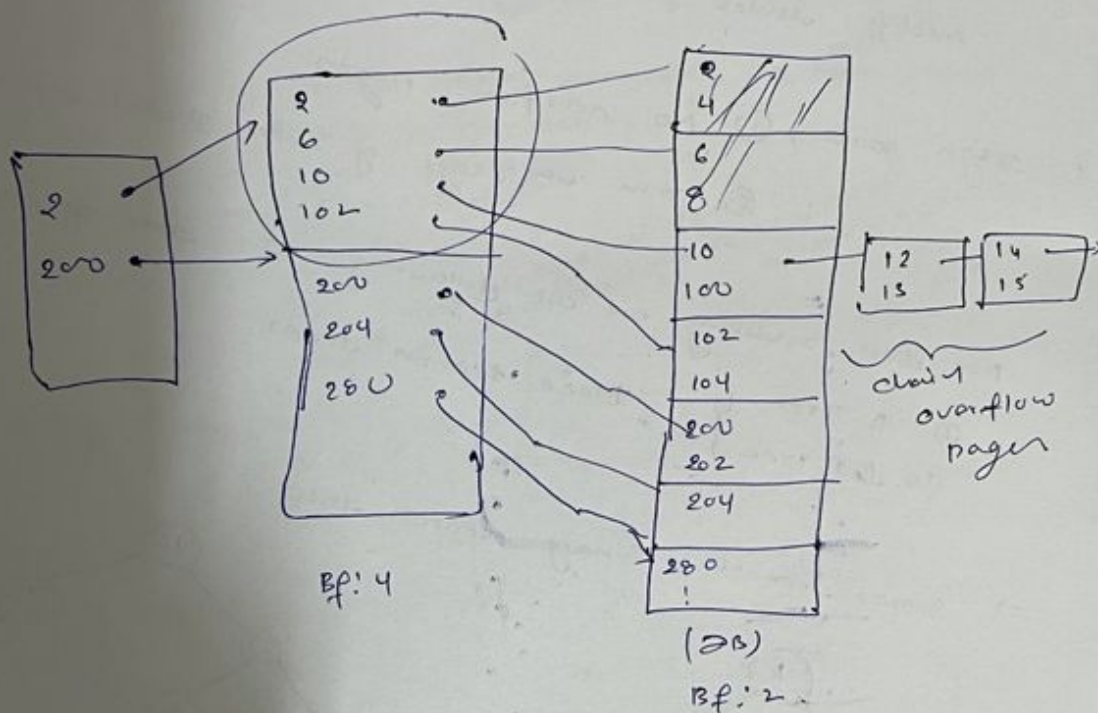
(a) Ans. $k = 3$ level

(b) total Block = $589 + 12 + 1$
= 602

(3) I/O cost = $k + 1$
= $3 + 1$
= 4 Block.

* Limitations of static Multilevel Index =
(as Indexing discussed before)

0% used.



* Insertion \Rightarrow overflow pages
 \downarrow
 worst case I/O cost $O(m)$
 m : # of DB blocks

* Deletion \Rightarrow minimum usage of Index Block is 1/2.

Because of these limitations we go through dynamic multilevel index.

* Dynamic Multi-level Index! - Based on Insertion/deletion of data record need to modify index file.

* Design goals \Rightarrow (1) No overflow pages
(2) min. use of index should be 50%

For this indexing we use following two data structures.

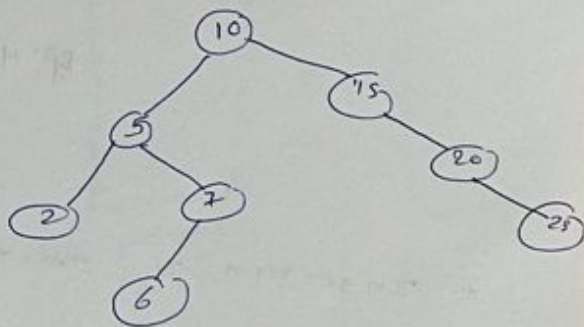
- (i) B-Tree
- (ii) B+ Tree

B+ tree search trees.

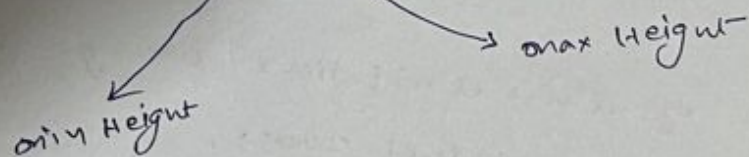
\Rightarrow Search tree \Rightarrow Binary search tree:



ET



\Rightarrow Binary search tree (N-keys)



Query processing overview

It is a procedure of transforming high-level query to low level language.

The basic steps of query ~~process~~ processing is ~~are~~ as follows:

→ Parsing & translation

→ Optimization

→ Evaluation

