

Question → Which cache line should be evicted?

Hawkeye {
→ Reconstructs belady's optimal for fast
→ Learn OPT behaviour of fast to inform eviction decisions for future loads

Two problems with OPT

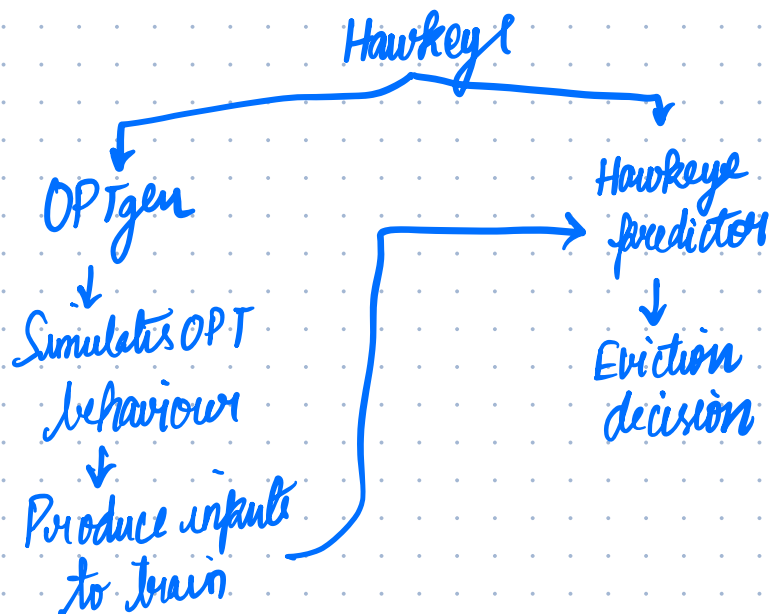
Efficient mechanism of reconstructing OPT → Solutions
→ Long history is needed to compute OPT → Set clumping

Key points: → OPT decision depends on the overlap of reuse intervals

→ OPT decision for fast access can be determined at next use

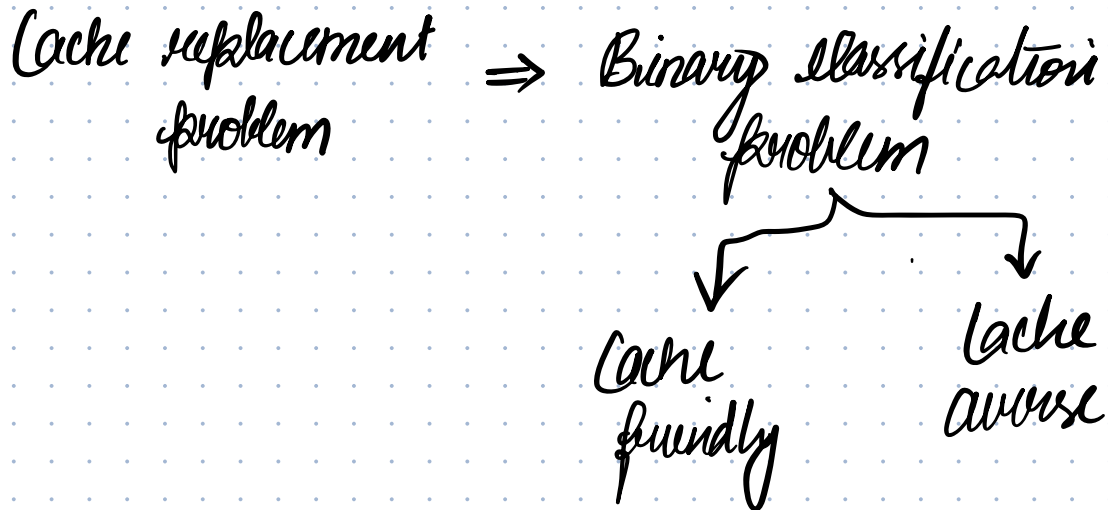
→ reuse of 8x is enough

Incoming line {
→ cache-friendly ⇒ Inserted with high priority
→ cache-averse ⇒ Marked as eviction candidates



Belady's OPT:

→ Evict the farthest reuse block.



* In fact if a certain load instruction has brought in lines which caused hits, then there is a high prob the same instruction would bring in lines that would be hit.

Doubt! PC ? \rightarrow Program counter?

\rightarrow Too much store overhead!

\downarrow
Set Dealing!

Hawkeye \rightarrow History \Rightarrow Sampled cache
 \rightarrow Decision \Rightarrow Occupancy vector & OPTgen
 \rightarrow Training \Rightarrow PC based predictor

Access



Sampled
cache

\rightarrow If tag match \rightarrow occupancy
vector



Binary
prediction



Train predictor



Assign
RRPV
value

Based
on
the PC



Cache-averse

Hit

1

Miss

1

Cache-friendly

0

0

age by 1
if all < 6 .

Optgen and occupancy vector:

- When a line is accessed put that as 0 in occupancy vector
- Check all past, increment that to this by 1. If exceeds Cache size, its a miss.

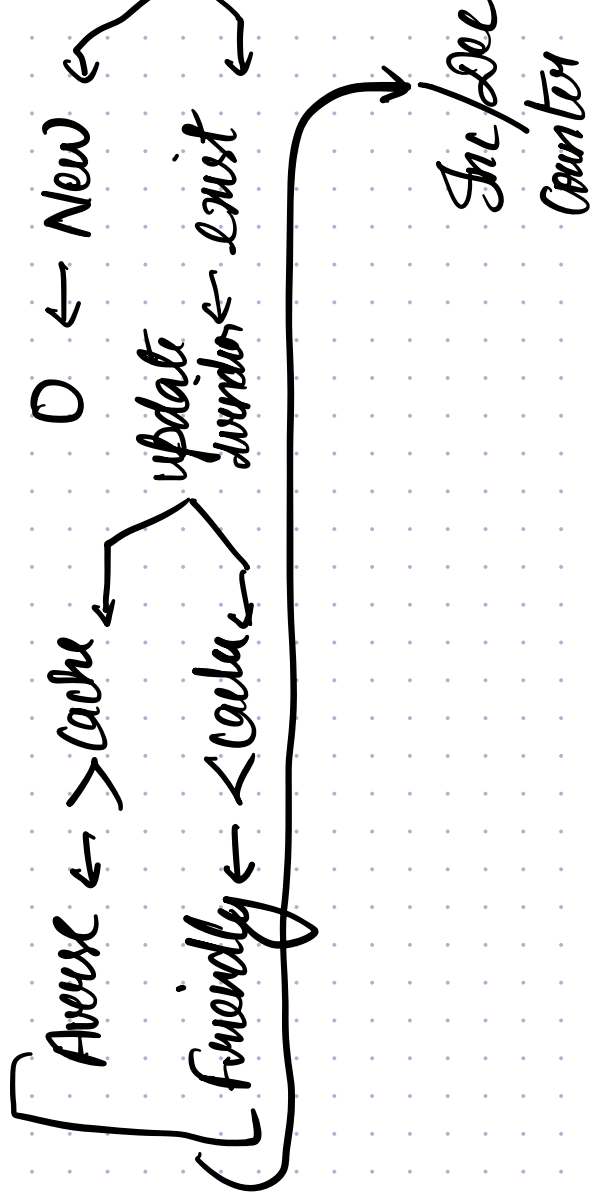
Doubt: Phase change boundary?

Evict cache aware

↓
No predicted Cache aware,
↓
LRU
↓
evicts cache-aware of new working set.

Prediction
detained
when ↓
Cache-friendly
lines are
evicted

Set-dueling



No averse

↓
Indicates
phase change

↓
LRU eviction

↓
detrain PC